Abstract

This Project focus on the scheduling problem. The simulation implements FCFS, SJF, RR and Priority methods and output corresponding output include average waiting time, average turnaround time, CPU utilization and Gantt chart. All the code is used C++ 11 standard and VS 2015 to compile.

1

The Shortest Job First scheduling algorithm is optimal in that it minimizes the average waiting time.

Proof

Denote there are n processes whose burst time is t_1 , t_2 , t_3 ... t_n , respectively. We also define that $t_1 < t_2 < t_3$... $< t_n$. Assume all processes arrive at the same time. The SJF algorithm results set is $\{n_1, n_2, n_3 ... n_n\}$, and average waiting time is

$$T = (t_1+t_2+...+t_n-1) / n$$

Assume SJF algorithm is not the minimum average waiting time algorithm. Thus, there must exist a set whose specific element n_j burst time larger than n_i where n_j 's burst time larger than n_i . And this makes $(T + t_j - t_i)/n < T$. However, $t_j > t_i$, which means the result is against the assumption. In other words, moving a short process before a long one decreases the waiting time of the short process more than it increases the waiting time of the long process. Consequently, T is the minimum time which means SJF is optimal in it minimizes the average waiting time.

2

2.1 Exercise 5.3

- a. The average turnaround time of FCFS algorithm is ((8+4-0.4)+(8+4+1-1)+8)/3 = 10.53
- b. The average turnaround time of SJF algorithm is (8+(8+1-1)+(8+4+1-0.4))/3 = 9.53
- c. The average turnaround time if CPU left 1 unit time and use SJF is ((1+1-1)+(4+1+1-0.4)+(8+4+1+1))/3 = 6.87

The Gantt charts show as follows:

FCF				
	P1	P2	F	23
0		8	12	13

SJF				
	P1	P3	P2	
0		8	9	13

SJF Delay

_				
		Р3	P2	P1
O	1	1 2	2	6 14

Exercise 5.12

The turnaround time and waiting time of each process show as table 2.1 and 2.2.

Table 2.1

_			1		
пu	rna	rou	na	l time	

	P1	P2	P3	P4	P5
FCFS	10	11	13	14	19
SJF	19	1	4	2	9
Nonpreemptive	16	1	18	19	6
Nonpreemptive Priority					
RR	19	2	7	4	14

Table 2.2

Waiting time

	P1	P2	P3	P4	P5
FCFS	0	10	11	13	14
SJF	9	0	2	1	4
Nonpreemptive	6	0	16	18	1
Nonpreemptive Priority					
RR	9	1	5	3	9

a. Average Turnaround time

FCFS
$$(10+11+13+14+19)/5 = 13.4$$

SJF
$$(1+2+4+9+19)/5 = 7$$

Priority (1+6+16+18+19)/5 = 12

RR (2+4+7+14+19)/5 = 9.2

b. Average Waiting Time

FCFS
$$(10+11+13+14)/5 = 9.6$$

SJF (1+2+4+9)/5 = 3.2

Priority (1+6+16+18)/5 = 8.2

RR (1+3+5+9+9)/5 = 5.4

c. The minimum average waiting time algorithm is Shortest Job First scheduling whose average waiting time is 3.2.

The four Gantt charts is showed as follows:

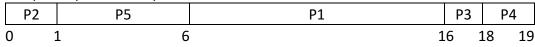
FCFS

. 0. 0								
	P1		P2	Р3		P4	P5	
0		10) 1	1	13	14		19

SJF

31									
P2	P4	Р3	P5	P1					

Nonpreemptive Priority



RR Ρ1 Р5 Р3 Р1 P2 Р3 Ρ4 P1 P5 Ρ1 P5 P1 P5 P5 P1 2 3 4 5 6 7 8 1 9 10 11 12 13 14 19 0

2.2 Simulation

The driven code of FCFS shows as figure 2.1, 2.2.

Figure 2.1

Figure 2.2

The SJF driven code shows as figure 2.3, 2.4, 2.5.

Figure 2.3

```
if (event_queue.empty())
{
    Event new_event = sjf_ready_queue.top();
    sjf_ready_queue.pop();
    event_queue.push(new_event);
    while (!sjf_ready_queue.empty()) {
        Event temp = sjf_ready_queue.top();
        if (temp.get_time()) = system_time + new_event.get_running_time())
        temp.update_time(system_time);
        if (ready_queue.pop();
        sjf_ready_queue.pop();
        sjf_ready_queue.pop();
        sjf_ready_queue.pop();
        sjf_ready_queue.pop();
        sstem_time += event.get_running_time(); // complete the process execution
        you stem_time += event.get_running_time(); // complete the process execution
        you stem_time += event.get_running_time() = e) { // record the waiting_time and turnaround time
        waiting_time[event.get_ID()] = e);
        turnaround_time[event.get_ID()] = event.get_running_time() - event.get_arrival_time();
        turnaround_time[event.get_ID()] = system_time - event.get_arrival_time();
```

Figure 2.4

Figure 2.5

The RR's driven code shows as figure 2.6, 2.7.

Figure 2.6

Figure 2.7

The Nonpreemptive Priority driven code shows as figure 2.8, 2.9.

Figure 2.8

Figure 2.9

The 5.3 and 5.12 output shows as figure $2.10 \sim 2.13$.

```
create PCB 1successfully
push event P2 into ready queue
create PCB 3successfully
push event P3 into ready queue
finish P1at time8
finish P2at time12
finish P3at time13
13
The Average waiting time of FCFS is 6.2
The Average waiting time of FCFS is 18.5333
! P1 ______ 8 ! P2 ____ 12 ! P3___ 13 !

create PCB 1successfully
push event P3 into ready queue
create PCB 2successfully
push event P2 into ready queue
create PCB 3successfully
push event P3 into ready queue
finish P1at time8
finish P1at time9
finish P2at time13
13
The Average waiting time of SJF is 5.2
The Average turnaround time of SJF is 9.53333
! P1 ______ 8 ! P2 ____ 13 !

create PCB 1successfully
push event P2 into ready queue
finish P3at time9
finish P3at time of SJF is 9.53333
! P1 _____ 8 ! P3__ 9 ! P2____ 13 !

create PCB 1successfully
push event P2 into ready queue
create PCB 3successfully
push event P3 into ready queue
fraish P3at time2
finish P3at time2
finish P3at time2
finish P3at time4

14
The Average vaiting time of delay_SJF is 2.86667
The Average vaiting time of delay_SJF is 6.86667
! delay_1 ! P3_2 ! P2___ 6 ! P1_____ 14 !
```

Figure 2.10

```
push event P3 into ready queue
create PCB 4successfully
push event P4 into ready queue
create PCB 5successfully
push event P5 into ready queue
finish P1at time10
finish P2at time11
finish P3at time13
finish P4at time14
finish P5at time17

19
The fiverage waiting time of FCFS is 9.6
The Reverage turnaround time of FCFS is 13.4
P1 10 P2 11 P3 13 P4 14 P5 19

create PCB 1successfully
create PCB 1successfully
push event P2 into ready queue
create PCB 3successfully
push event P3 into ready queue
create PCB 4successfully
push event P4 into ready queue
create PCB 4successfully
push event P5 into ready queue
create PCB 4successfully
push event P5 into ready queue
finish P3at time4
finish P3at time4
finish P3at time4
finish P3at time6
finish P3at time6
finish P3at time7

19
The Reverage waiting time of SJF is 7
P2.1 P4.2 P3.4 P5.9 P1. 19
The Reverage turnaround time of SJF is 7
P2.1 P4.2 P3.4 P5.9 P1. 19
The Reverage turnaround time of SJF is 7
P2.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
P2.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
P2.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
P3.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
P3.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
P3.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
P3.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
P4.1 P4.2 P3.4 P5. 9 P1. 19
The Reverage turnaround time of SJF is 7
```

Figure 2.11

```
push event 2 into ready queue
create PCB Successfully
3 running tine is 2
push event 2 into ready queue
create PCB Successfully
4 running tine is 2
push event 3 into ready queue
create PCB Successfully
4 running tine is 1
push event 4 into ready queue
push event 5 into ready queue
push event 6 into 8 int
```

Figure 2.12

Figure 2.13

3

The simulation uses random exponential generator to generate the I/O interrupt randomly.

a) Observation of FCFS and SJF

Because CPU does not know the processes' burst time. Thus, CPU will predict the CPU burst time to run SJF methods. The equation is as follows:

$$\tau_{n+1} = \alpha t_n + (1-\alpha)\tau_n \qquad (0 \le \alpha \le 1) \text{ in }$$

The results show as table 3.1.

Table 3.1 average waiting time

	Average	Average	Throughput	CPU	Total
	waiting	turnaround	(hour)	utilization	simulation
	time(min)	time(min)			time(min)
FCFS	24.3727	27.1323	17.8685	0.821821	33.578667
SJF(α=1)	24.3165	27.076	17.9139	0.823907	33.493667
SJF(α=0.5)	24.3354	27.0949	17.8107	0.819162	33.687667
SJF(α=1/3)	24.2546	27.0142	17.8451	0.820746	33.622667

From table 3.1, when α is 0.5, the system has the lowest average waiting time, however, the simulation time is the longest. And all SJF simulations' average time is

smaller than FCFS but the FCFS simulation has a good CPU utilization compared with SJF α is 0.5 and α is 1/3. The average waiting time is good, but the CPU utilization is a little different because the random generator makes the IO switching times more or less to results the FCFS has a good CPU utilization.

b) Statement:

- 1) When quantum is small enough, the CPU utilization time will increase when quantum increase. And the average waiting time will decrease when quantum increase.
- 2) With increasing the quantum close to the IO requests time, the CPU utilization and average waiting time will vibrate follow the time.
- c) Based on the statement at b), the system results show as figure $3.1 \sim 3.6$.

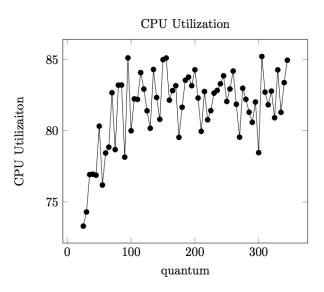


Figure 3.1

From figure 3.1, we know that the quantum from 0 to 50, the CPU utilization is increasing. After 50, the CPU utilization has a vibrate data. It's because when quantum is small enough (i.e. 25), the system timer interruption is almost determined by quantum. And the quantum is too small make the whole system switching too many times which results low CPU utilization. With the quantum increasing to the value which the timer interruption determined not only by quantum, but also IO requests, the CPU utilization determined by the interrupt time which is better to perform the CPU utilization.

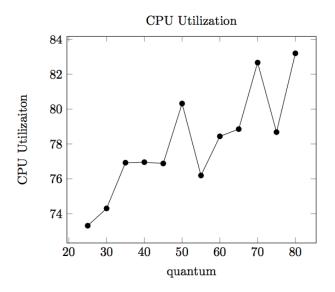


Figure 3.2

From the figure 3.2, the conclusion is that the CPU utilization's trend is increasing with the quantum from 20 to 80 because the the system switching times reduces. Every process has more time to run in every single timer.

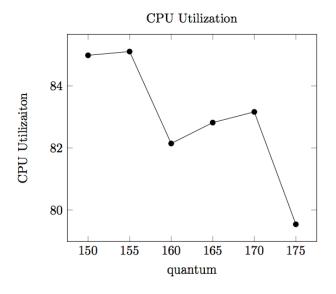


Figure 3.3

From figure 3.3, the CPU utilization decreases with quantum from 155 to 175. Because of the timer interruption decreases, the system efficiency is growing down. Consequently, the timer interruption length is important to the system efficiency. If it is too short, the system will have too many switching times, which results CPU utilization low. If it is too long, the system will have bad performance like FCFS. The system CPU utilization is vibrating because it the IO requests is random generated and this is not stable.

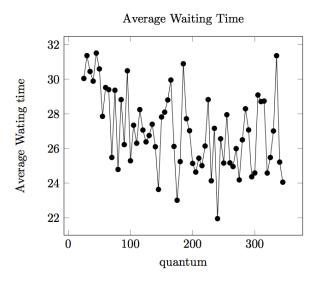


Figure 3.4

From figure 3.4, the average waiting time tendency is decrease with quantum from 0 to 80. Because during this quantum values, the system switching times decreases and increases the CPU utilization which makes the system is more efficiency.

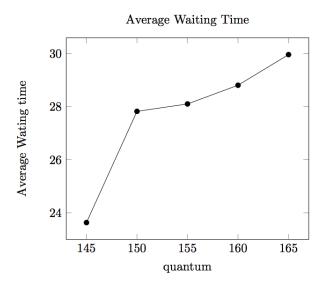


Figure 3.5

From figure 3.5, the average waiting time increase because the quantum becomes large and make the system degenerate to FCFS. Thus, the system average waiting time is larger and larger.

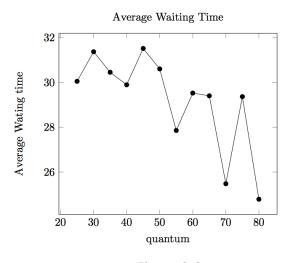


Figure 3.6

From figure 3.6, the average waiting time has a decrease trend because the with the quantum increasing the context switching times decreases in a reasonable scale. And this will make system reduce large amount of switching time and waiting time.