It Problem No: 03

Ptroblem name: Suppose that the relation R1 and R2 on a set A are trepresented by the matrices $M_{R1} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \text{ and } M_{R2} = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \end{bmatrix}. \text{ Write a program to find the MR1 UR2 and MR1 <math>\Theta$ R2.

Theorey: The preogram shown above defines two function.

I matrix union () and conatrix exclusive ore (), to

compute the union and symmetric curion (exclusive ore)

of two matrices respectively.

The matrix union of function takes two mostrics as input and treturens a new mostrics that represented the union of two input matrices. It uses list competitions to perform element-wise logical or operation between the element of the two input mostrices.

The motrik-exclusive-orc () function also takes two motrics as input and reducting a new motrics that represent the symmetric difference of the two input matrics. It uses list competitions and the X-or operators (1) to percotor element wise logical X-or operation between the elements of the two input matrics.

After defining this functions, the preogram eneades two motrics R1 and R2 which represent relations. These motrics are then passed to the motrics unions and matrix exclusive orz () functions to compute the union and symmetric difference MR1 PR1.

抽 Here a python program is written below:

def matrix union (matrix 1, matrix 2):

union-tresult = [[matrix1[i][j] or matrix2[i][j]

fon j in range (len (matrix1[o]))]

fon i in trang (len (matrix1))]

return-union-tresult

Function to compute the motrix union (element-wise or)

def mattix_elexalusive_or (matrix 1, matrix 2):

XOR result = [[matrix][i][j] ^ matrix2[i][j]

For j in trange (len (mostrix1[6]))] for i in trange (len (mostrix))]
tratation xor result.

Function to compute the mostrix exclusive or (elementwise XOR) of the two mostrix.

R1=
$$[[1,0,1]$$
, # Give the monthix representing redation $[1,0,0]$, R1 and R2. $[0,1,1]$

$$R2 = \begin{bmatrix} [1,0,1] \\ [0,1,1] \\ [2,0,1] \end{bmatrix}$$

MRIO

mR1xorr_R2 = moditix_exclusive_orr (R1, R2) # find mR10R2

mR1_union_R2 = moditix_union (R1, R2) and MR1 UR2.

print ("Firest Matrix = ", R1)

prant ("Second Matrix = " R2)

print ("Rows =" len (R1), "Co15 = ", len (R1 [0]))

preint ("Matrix-union RRMRIUMR2", MRI-union-R2)

Ptcint (" Matrix Exclusive or mR1 AR2", mR1 DR2)

preint the result,

Output:

Firest Matrix = [[1,0,1],[1,0,0],[0,1,1]]

Second modraix = [[1,0,1], [0,1,1], [1,0,1]]

Rows = 3 Cols = 3

Modreix Union mRIUR2: [[1.0.1], [1.1.1], [1.1.1]

Matrix Exclusive OR MRA @RL: [[0,0,0], (1,1,1], [[,1,0])

X

A Problem no - y

Parts by Wateshall's algorithm.

Theorey & the program shown above applies the floyd-Warrshall algorithm to find the shortlest paths in a greath. The graph is represented by an adjacency matrix.

The function 'floydwatshall()' takes two algorith: The numbers of vertices in the grouph and the adjacency modicing nepresenting the grouph. Within this function, the floyd wareshall algorithm is implimented using nested forcloops

After applying the floy of workshall algorithm, the progressing prints the resulting adjacency mostrix, showing the distance between each pain of veretices. The left most trom represents the origin veretex and the topmost column represents the destination veretex. The values in the mostrix represent the shortest distance between the cornesponding veretex.

In the main () function the number of vertices is set to 4. The adjancency matrix is defined best based on the given graph.

Finally the mains' function calls the floy distribull' function with the number of vertices and the adjacency matrix as areguments to find the shoutest paths in the grouph.

Here a C++ ptrogtrom is written given below:

Include (iostream)
using namespace std; // Defining the infinity value,
Const int INF = 1000000000;

// Function to apply floyd-Warshall algorithm
void floydwarshall (int veretex, int adjancency_matrix[][4])

// I-letrate over all vetetex as intermediate notes
for (int k=0: K< veretex; K++)

Ill-for each pair of vertex (i,i), check if going to vortex k provides a short path for (int i = 0; i < vertex; i++)

P-T.O.

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Ten (intj=0; j\vertetex; j++)
   In Relux the distance from i to j by allowing ventex k as an
    intermediate ventex
   Il consider which one is better going through ventex k on
    the previous value
    adjacenez-matrix[i][j]=min(adjancenez-matrix[i][f]
           adjancency_motrix [i][k] + adjacency_motrix [k][j]);
11 Pretty print the graph
11 0/d means the leftmost trow is the origin veretex.
Il and the topmost column as distinction verotex
 cout<< " 0/d" >
for (int i=0; i< vertex; i++)
 ? cout<< 't'<< i+1;
 cout << end 1;
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U
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Il pranting the adjordency montrux
for (int i= 0; i(vetctex; i++)
} cout << i+1;
 for (int j = 0; j < vertex; j++)
 ) cout<<'t'<< adjacency_matrix [i][j];
 aout (cend L;
int main ()
 1/ Number of verties in the graph
 int vertex = 4;
Input is given as adjacency mostrix,
input represent this our directed grouph
Should set infinite value for each palor of vertex
 fact has no edge
```

11 The adjacency mortrix represented fre graph int adjacency_mostrix [4][4] = ? 30,5, INF, 103, ? INF, 0, 3, INF }, ? INF, INF, 0, 1}, 3 INF, INF, INF, 0}); If find the shortest posts using Floy & wareshall algorithm by call the function floydwarshall (vettex, adjacency_matrix); réfusino; Output 8 4 1 0 5 2 1000000000 ધ 1000000000 3 1000000000 100000000 1000000000 100000000