

# Project 1: Mathematical and Empirical Analysis

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1. Pseudocode describing your reformat date algorithm.

define global monthArr and initialize with full name of Months

define global monArr and initialize with first three letters of months

// validate range helper function

bool validateRange(int year, int, month, int day):

    check if yer between 1900 and 2099 else throw exception

    check if month between 1 and 12 else throw exception

    check if day between 1 and 31 else throw exception

return true

// is\_digits function to verify if numbers

bool is\_digits(string& str):

    return true if all are digits else return false

// check pattern function

int checkPattern(input):

    if not is\_digits(input) : throw exception

    check if input contains " - ", if so return 1

    check if input contains " / ", if so return 2

    check if input month is full month name, if so return 3

    check if input month is first 3 characters of month, if so return 4

throw exception otherwise

def reformat\_date(string input) :

    pattern = checkPatter(input) (valid patterns: 1, 2, 3, 4)

    check if pattern is valid : otherwise throw invalid\_argument

    // needed variables

    output set to empty string

    year set to 0

    month set to 0

    strMonth set to empty string

    strDay set to empty string

    If (pattern == 1):

        make temp index

        assign values to year, month, and day

        Increment temp index to account for '-'

        if (validateRange(year, month, day):

            check that year, month, day are within bounds set

            concatenate to output variable and return

    If (pattern == 2):

        make temp index

        assign values to year, month, and day

        Increment temp index to account for '/'

        if (validateRange(year, month, day):

            check that year, month, day are within bounds set

            concatenate to output variable and return

```
if (patter == 3) :  
    make a temp index  
    assign month to inputMonth variable  
    increment temp index  
    for loop to check if inputMonth is in monthArr  
    assign value to day  
    increment temp index based on comma  
    assign value to year  
    if (validateRange(year, month, day):  
        check that year, month, day are within bounds set  
if (pattern == 4):  
    make a temp index  
    assign month to inputMonth variable  
    increment temp index  
    for loop to check if inputMonth is in monArr  
    assign value to day  
    increment temp index based on a space character  
    assign value to year  
    if (validateRange(year, month, day):  
        check that year, month, day are within bounds set  
    concatenate to output variable and return
```

## 2. Mathematical analyses for each of the three algorithms.

```

run_length_encode(S):
    C = "" 1
    if S is empty: 1
        return C 1
    run_char = S[0] 1
    run_length = 1 1
    for each character c in S after the first: n
        if c == run_char: n
            run_length++ n
        else:
            append_run(C, run_char, run_length)n
            run_char = c n
            run_length = 1 n
            append_run(C, run_char, run_length) 1
    return C 1

append_run(C, run_char, run_length): 1
    if run_length > 1: 1
        C.append(convert_to_string(run_length)) 1
        C.append(run_char) 1

```

$$T(n) = 5n + 11$$

By properties of O

$$5n + 11 \in O(5n + 11) \quad (\text{trivial})$$

$$= O(\max(5n, 11)) \quad (\text{Dominated terms})$$

$$= O(5n)$$

$$= O(n) \quad (\text{Constant factor})$$

```

longest_frequent_substring(S, k):
    freq = new dictionary that maps a char to an int      1
    for c in S:                                           n
        if c not in freq:                                n
            freq[c] = 1                                   n
        else:
            freq[c]++                                     n
    best = ""                                             1
    for b from 0 through n-1:                             n
        for e from b+1 through n:                         n*n
            cand = S[b:e]                                 n*n*n
            if every character c in cand has freq[c] >= k: n*n*n
                if cand.size() > best.size():             n*n*n
                    best = cand                           n*n*n
    return best                                           1

```

$$T(n) = 4n^3 + n^2 + 5n + 3$$

By properties of O

$$4n^3 + n^2 + 5n + 3 \in O(4n^3 + n^2 + 5n + 3) \quad (\text{trivial})$$

$$= O(\max(4n^3, 5n^2, 5n, 3)) \quad (\text{Dominated terms})$$

$$= O(4n^3)$$

$$= O(n^3) \quad (\text{Constant factor})$$

## Reformat Date Algorithm

$$T(n) = 2n^2 + 9n + 8$$

By properties of O

$$2n^2 + 9n + 8 \in O(2n^2 + 9n + 8) \quad (\text{trivial})$$

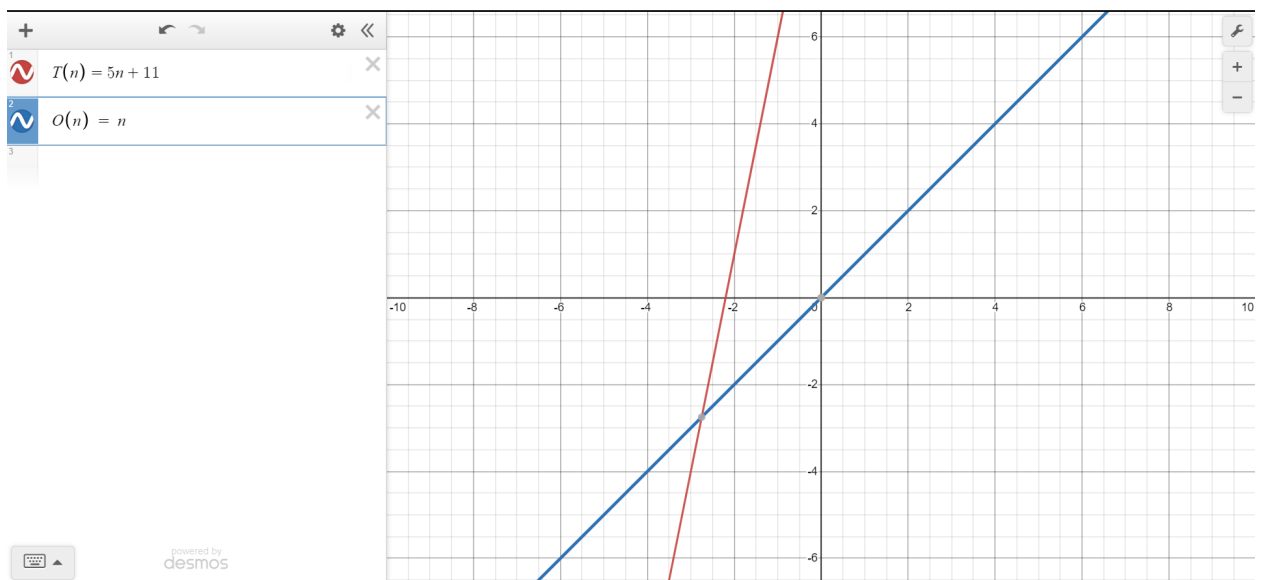
$$= O(\max(2n^2, 9n, 8)) \quad (\text{Dominated terms})$$

$$= O(2n^2)$$

$$= O(n^2) \quad (\text{Constant factor})$$

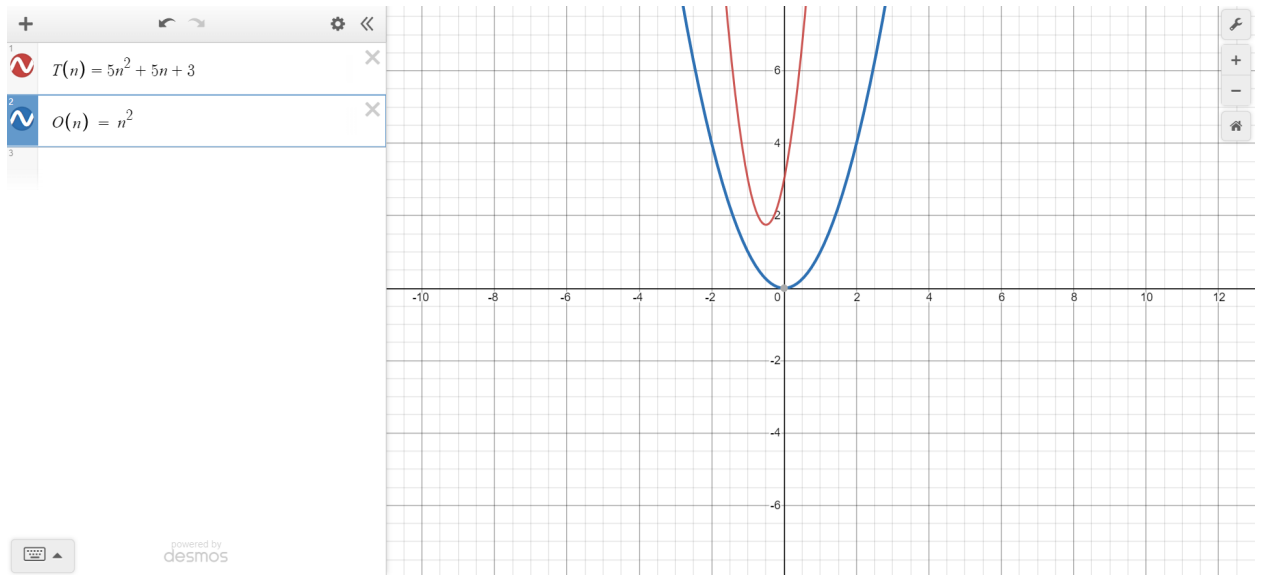
3. Scatter plots for each of the three algorithms.

Algorithm 1

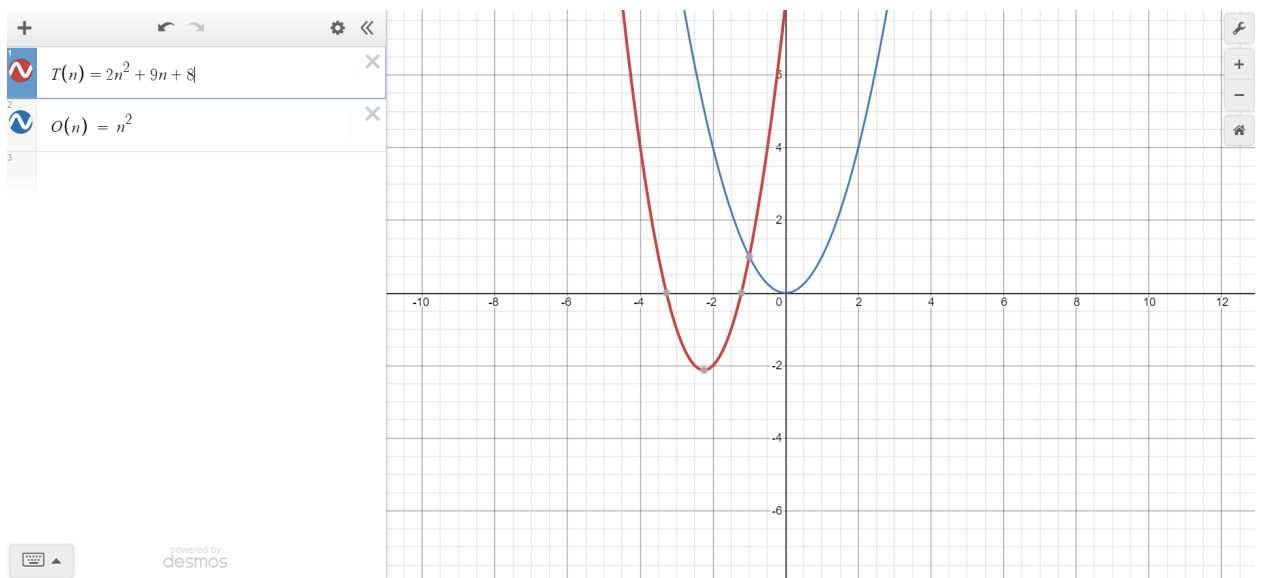


Algorithm 2





### Algorithm 3



4. Answers to the following questions. (Each answer should be at least one complete sentence.)
1. What is the efficiency class of each of the algorithms, according to your own mathematical analysis? (You are not required to include all your math work, just state the classes you derived and proved.)

The efficiency class for each of the algorithms according to our mathematical analysis and by properties of  $O$  came out to be for run-length encode  $O(n)$ , longest frequent substring as  $O(n^2)$ , and for the third algorithm  $O(n^2)$ .

2. Between the run-length encode and longest frequent substring algorithms, is there a noticeable difference in the running speed? Which is faster, and by how much? Does this surprise you?

After running the make test command a few times it looks like the Longest frequent substring is much faster. Last run for run-length encode was 1098ms vs 31ms for Lfs. This is surprising since Lfs has a nested for loop.

3. Are the fit lines on your scatter plots consistent with the efficiency classes predicted by your math analyses? Justify your answer.

The fit lines on the scatter plot are consistent with efficiency classes that were predicted by the math analysis because the shape of  $T(n)$  and Big  $O$  are generally the same shape. However, the graphs generated from  $T(n)$  are usually displaced on the x-axis or y-axis and can sometimes be narrow or wider compared to the graph generated from the Big  $O$ .

4. Is all this evidence consistent or inconsistent with the hypothesis stated on the first page? Justify your answer.

From the looks of all the evidence, for large values of  $n$  the derived efficiency classes were accurately predicted for the observed running time of the implementation.