

# **Conflict-free Replicated Data Type (CRDT)**

# What is a Lock-Free Data Structure?

A lock-free data structure enables concurrent, non-blocking access to shared data without traditional locks, boosting performance in multi-threaded applications.

## Key Concepts

- **No Waiting:** Avoids lock contention to enhance performance.
- **Atomic Operations:** Uses low-level atomic operations (e.g., Compare-And-Swap) for safe, consistent data updates.
- **High Concurrency:** Improves performance by allowing multiple threads to interact without blocking.

## Analogy

Like updating a shared Google Doc — everyone edits freely, last valid edit wins.

# LOCK FREE DATA STRUCTURES

## PROS

Enhanced Performance: Threads operate concurrently without locks, boosting throughput and performance in contested scenarios.

No Deadlocks: Lock absence eliminates deadlocks, increasing system robustness and reliability under heavy loads.

## CONS

Implementation Complexity: Designing lock-free algorithms is exceptionally difficult, requiring expert knowledge to manage thread synchronization and prevent race conditions.

Hardware Dependency: Lock-free algorithms lack universal portability, relying on CPU-specific atomic instructions (e.g., CAS) that necessitate code adaptation for different hardware architectures.

# Message Passing Model of Process Communication

This model allows processes to communicate without sharing the same address space. They exchange information by sending messages to and receiving messages from a common **message queue**.

- Each thread (actor) has a **mailbox**
- Avoids shared-state issues

## Why Message Passing Helps?

- No race conditions on shared memory
- Easier reasoning about state
- Fits distributed systems naturally

# Broadcasting Local Updates via Message Passing

- **User Makes Local Edit** → Change detected
- **System Creates Update Object** → captures *what changed, where, when, by whom*
- **Query Shared Registry** → find all active users & their message queues

# Conflict-Free Replicated Data Type (CRDT)

A CRDT is a special type of data structure designed to handle concurrent updates across multiple devices or servers (**replicas**) without the need for a central coordinator.

## Use Cases:

- **Mobile Apps:** Syncing data (notes, calendars) across a user's devices.
- **Distributed Databases:** Ensuring data consistency across multiple database nodes.
- **Collaboration Software:** Allowing multiple users to edit the same document simultaneously.

## Types of CRDTs:

There are two main types of CRDTs:

- **State-based CRDTs:** These replicate the entire state of the data structure. Merging is done by comparing and reconciling the states of different replicas.
- **Operation-based CRDTs:** These replicate the operations performed on the data structure. Merging is done by applying operations from different replicas in a commutative manner.

# HOW CRDTs work?

All operations follow these simple rules:

- **Commutativity:** Order doesn't matter.

$$(\text{Update A} + \text{Update B}) = (\text{Update B} + \text{Update A})$$

- **Associativity:** Grouping doesn't matter.

$$((A+B) + C) = (A + (B+C))$$

- **Idempotence:** Duplicates have no extra effect.

$$(\text{State} + \text{Update A} + \text{Update A}) = (\text{State} + \text{Update A})$$

# Questions?