



Complete - every level (Except possible last) is full. And last fills in from left.

Become a heap if every parent  $\geq$  Child

The reason we haven't talked about the reason this needs to be a complete binary tree is because you can store it in an array - tree structure is implicit

Heap Sort is  $O(n \lg n)$

If parent index is  $i$ , children indices are  $2i$  and  $2i+1$   
 Paren of node  $i$  are  $x = \text{floor}(i/2)$

Build a heap:  
 Bottom up  $O(n)$

First finds the index of the largest element that still has children.  
 In this case the index would be 6 with element 9. 9 is greater than all of its children so it's good

Now work backwards, check the 5 node. 5 is  $\geq$  than all of its children so it's good

1 is not good so 6 and 1 are switched

Keep doing this

Repeat:  
 Remove Largest element -  $O(1)$   
 Reform tree -  $O(\lg n)$