

REQUIREMENTS NOT MET

N/A

VIDEO FILE LINK

<https://youtube.com/shorts/TeFolGsbdYU?feature=share>

PROBLEMS ENCOUNTERED

I had very minor issues with my first attempts at my ASMs, but my PI helped me fix them. A major problem I had was creating the ROM and MIF files; I found the tutorials confusing, and it took me much longer than it should have. One issue I had was that I flipped the PC's MSB and LSB connections, which caused incorrect ROM access. The final issue I had was with creating the simulations, as it involved a lot of troubleshooting and minor corrections before they were correct.

FUTURE WORK/APPLICATIONS

The CPU designed in Lab 6 shows how a state machine can execute instructions using a custom instruction set. This basic architecture can be extended to support features like branching, conditional logic, and memory addressing. In Lab 7, these concepts are applied to the G-CPU, a more advanced processor capable of running full assembly programs. The skills developed in Lab 6 provide the foundation for writing, assembling, and simulating code in a complete CPU system.

PRE-LAB QUESTIONS OR EXERCISES

PART 1 PRE-LAB QUESTIONS

1. Why did we require the new instruction register in this design?

The instruction register is required to temporarily store the opcode of the current instruction, allowing the controller to decode it and generate the necessary control signals for the data path. This design allows the CPU to execute instructions in a synchronized and efficient manner across clock cycles, ensuring consistent operation.

2. In this section of the lab, you are setting the INPUT bus by hand. If you wanted to read or fetch this value from memory, what could you add to do this automatically for you every CLK cycle?

By adding a program counter and a ROM module, the system could automatically fetch opcodes and data from memory on each clock cycle. The PC would sequentially supply memory addresses, and the ROM would provide the stored instructions, eliminating the need for manual input.

3. How would you add more instructions (i.e., 8 instead of 4) to the controller?

To accommodate eight instructions, the instruction register would need to be expanded from 2 bits to 3 bits, allowing for eight unique opcodes. The controller's state machine logic would also need to be updated to decode the additional opcodes and produce the appropriate control signals for each operation.

PART 2 PRE-LAB QUESTIONS

1. Why do we need the extra states in the LDAA and JMP instruction paths?

The LDAA and JMP instructions require extra states because they each involve two separate memory reads, which must occur over two clock cycles. For LDAA, the first cycle fetches the opcode, and the second retrieves the immediate data to load into REGA. For JMP, the first receives the opcode, and the second cycle retrieves the target address to load into the Program Counter.

2. What do you need to do to the address lines to get your program to start at address \$2CD0 (instead of \$3A10)?

You would have to tie the upper address lines, from 14 to 4, to \$2CD, which is binary 0010 1100 1101. Addresses 4, 6, 7, 10, 11, and 13 would be tied to VCC, and the rest would be tied to ground.

PRE-LAB REQUIREMENTS (Design, Schematic, ASM Chart, ASM Chart, VHDL, etc.)

1. FIRST RALU CONTROLLER:

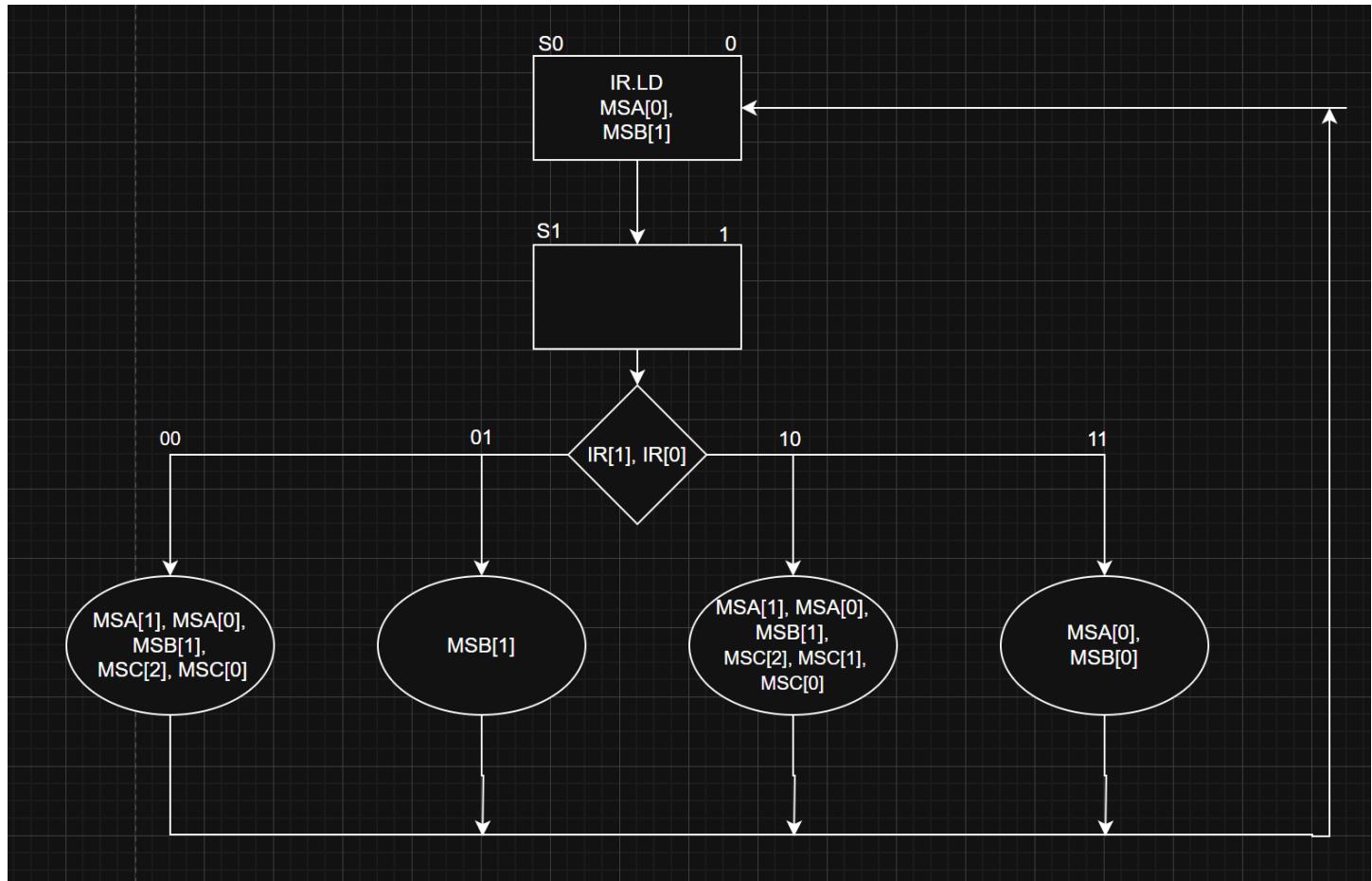
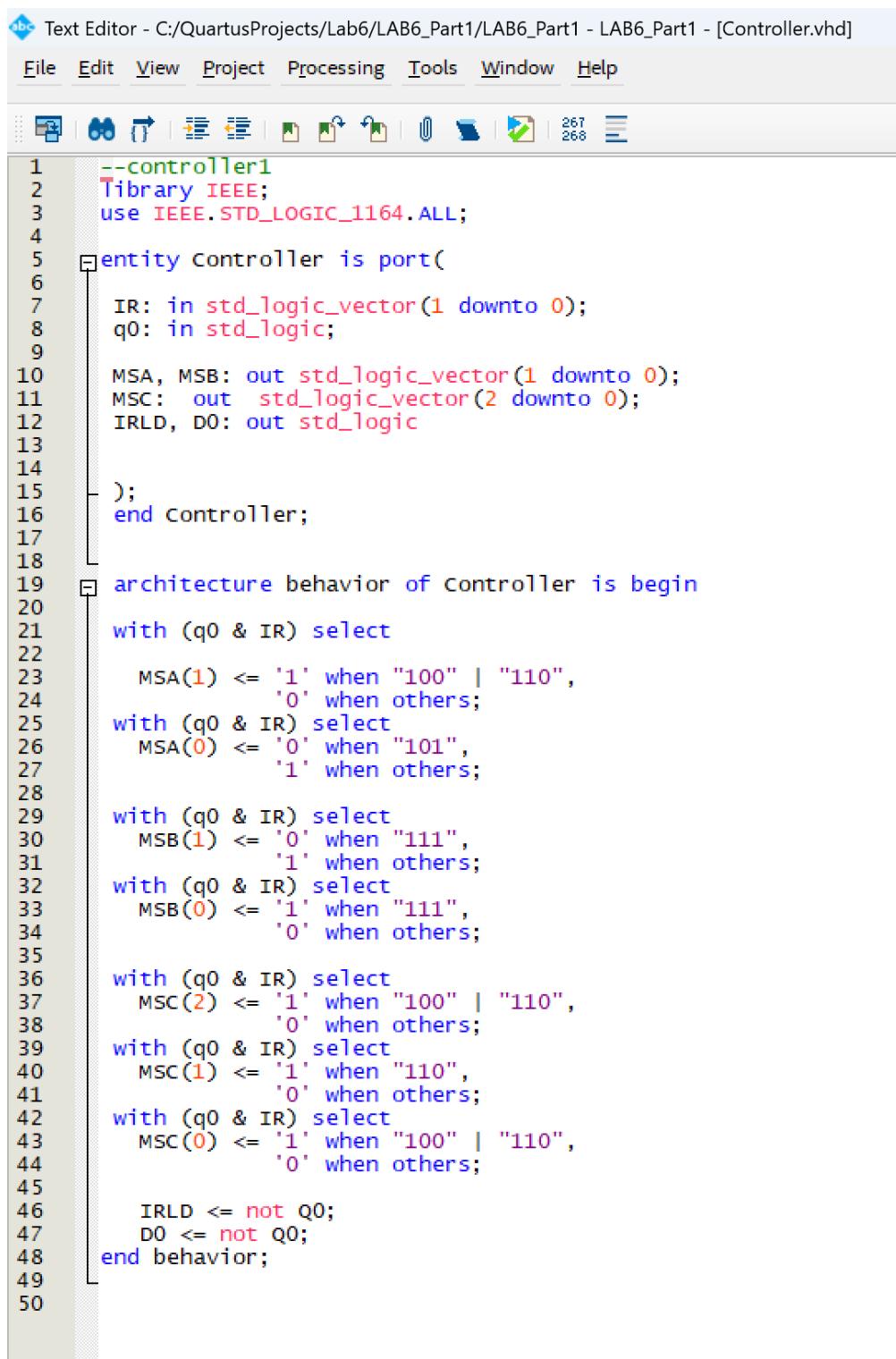


Figure 1: ASM for First RALU Controller (Part I)

Q0	IR[1]	IR[0]	IR.LD	MSA[1]	MSA[0]	MSB[1]	MSB[0]	MSC[2]	MSC[1]	MSC[0]	Q0+
0	-	-	1	0	1	1	0	0	0	0	1
1	0	0	0	1	1	1	0	1	0	1	0
1	0	1	0	0	0	1	0	0	0	0	0
1	1	0	0	1	1	1	0	1	1	1	0
1	1	1	0	0	1	0	1	0	0	0	0

Table 1: NSTT for First RALU Controller (Part I)

Lab 6 Report: Elementary CPU Design



The screenshot shows a Quartus II Text Editor window with the following details:

- Title Bar:** Text Editor - C:/QuartusProjects/Lab6/LAB6_Part1/LAB6_Part1 - LAB6_Part1 - [Controller.vhd]
- Menu Bar:** File Edit View Project Processing Tools Window Help
- Toolbar:** Includes icons for New, Open, Save, Undo, Redo, Find, Replace, Copy, Paste, Cut, Select All, and others.
- Code Area:** Displays VHDL code for a controller entity named "Controller". The code defines port connections for IR, q0, MSA, MSB, MSC, IRLD, and D0. It then defines an architecture "behavior" for the controller, which uses multiple nested selects based on q0 and IR values to assign values to MSA, MSB, and MSC outputs. It also sets IRLD and D0 based on q0.
- Line Numbers:** Numerical line numbers are displayed on the left side of the code area, ranging from 1 to 50.

Figure 2: VHDL for First RALU Controller (Part 1)

Lab 6 Report: Elementary CPU Design

Lab 6 Part 1 (FIRST RALU CONTROLLER)
 Name: Arion Stern
 Class #: 10844
 PI Name: Erick Zayas Ramos
 Description: Implementing a RALU Controller for the LAB4 RALU

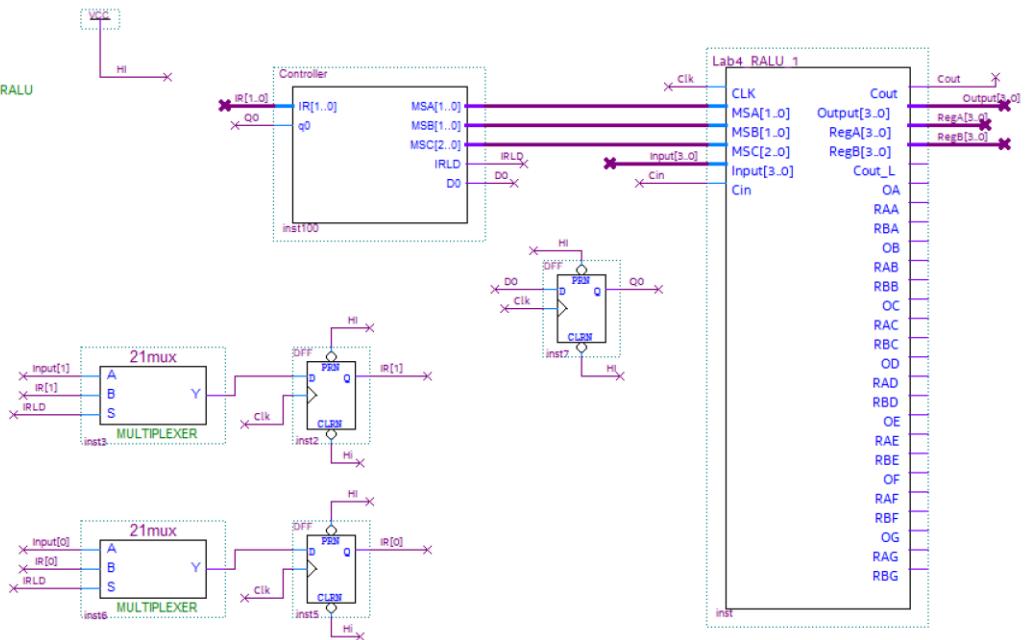


Figure 3: BDF Schematic for LAB6_Part1 (Part 1)

Whenever the controller in state 0 the IRLD is true and the two LSB input bits are loaded into IR at the clock edge.

When the controller is in state Q1 it performs the correct function according to the IR bits and the registers get loaded at the next clock edge.

Showing that Cin works with SUMAB:

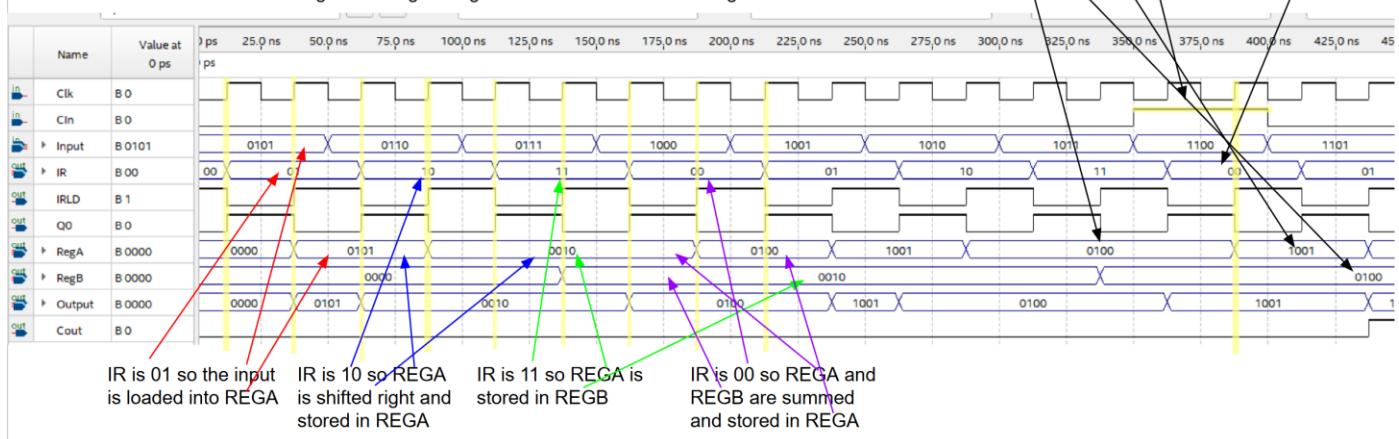


Figure 4: Annotated Simulation for Lab6_Part1 (Part 1)

Lab 6 Report: Elementary CPU Design

2. SECOND RALU CONTROLLER WITH ROM :

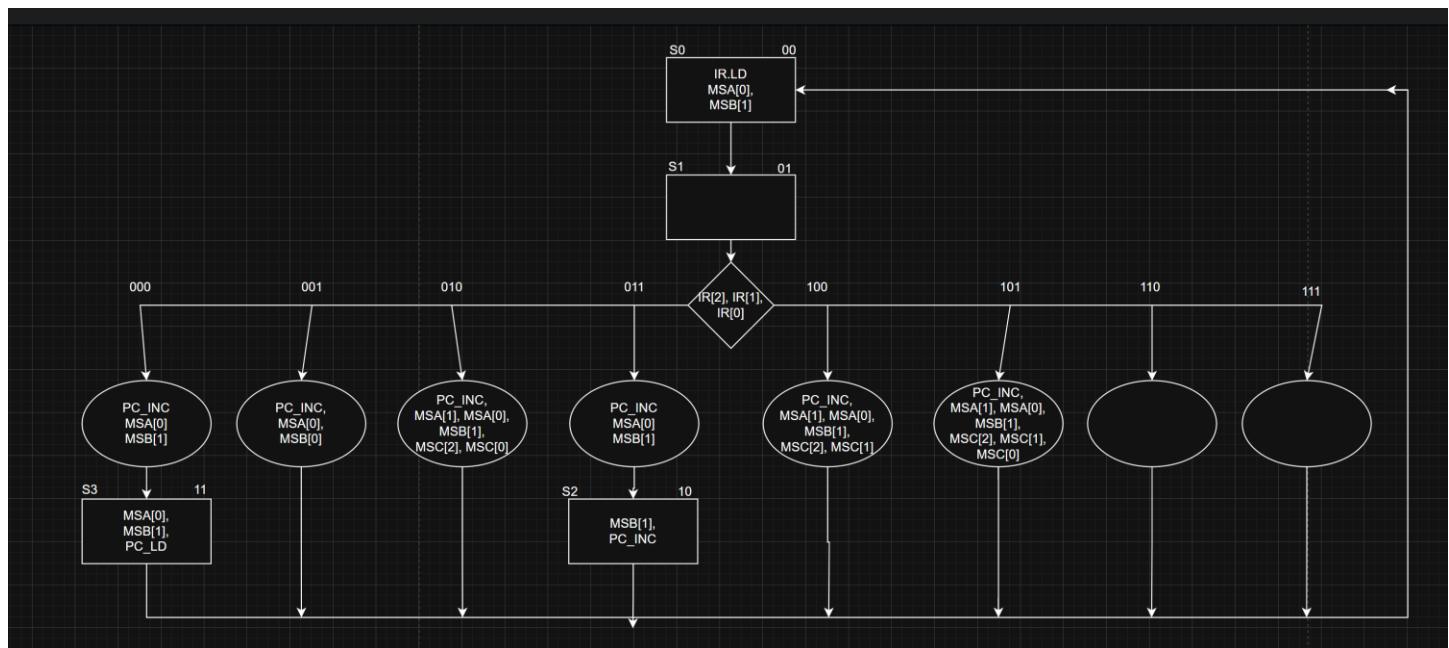


Figure 5: ASM for Second RALU Controller (Part 2)

Q1	Q0	IR[2]	IR[1]	IR[0]	IR.LD	PC.LD	PC.INC	MSA[1]	MSA[0]	MSB[1]	MSB[0]	MSC[2]	MSC[1]	MSC[0]	Q1+	Q0+
0	0	-	-	-	1	0	0	0	1	1	0	0	0	0	0	1
0	1	0	0	0	0	0	1	0	1	1	0	0	0	0	1	1
0	1	0	0	1	0	0	1	0	1	0	1	0	0	0	0	0
0	1	0	1	0	0	0	1	1	1	1	0	1	0	1	0	0
0	1	0	1	1	0	0	1	0	1	1	0	0	0	0	1	0
0	1	1	0	0	0	0	1	1	1	1	0	1	1	0	0	0
0	1	1	0	1	0	0	1	1	1	1	0	1	1	1	0	0
1	0	-	-	-	0	0	1	0	0	1	0	0	1	0	0	0
1	1	-	-	-	0	1	0	0	1	1	0	0	0	0	0	0

Table 2: NSTT for Second RALU Controller (Part 2)

Lab 6 Report: Elementary CPU Design

```
--controller2
library IEEE;
use IEEE.STD_LOGIC_1164.ALL;
entity controller2 is port (
    Q1, Q0: in std_logic;
    IR: in std_logic_vector(2 downto 0);
    IRLD, PCLD, PCINC, D1, D0: out std_logic;
    MSA, MSB: out std_logic_vector (1 downto 0);
    MSC: out std_logic_vector (2 downto 0)
); end controller2;
architecture behavior of controller2 is
    signal state : std_logic_vector(1 downto 0);
begin
    state <= Q1 & Q0;
    IRLD <= not (Q0 or Q1);
    PCLD <= q1 and q0;
    --pos 00 and 11
    PCINC <= (Q1 or Q0) and ((not q1) or (not q0));
    with (q1 & q0 & IR) select
        MSA(1) <= '1' when "01010" | "01100" | "01101",
        '0' when others;
    with state select
        MSA(0) <= '0' when "10",
        '1' when others;
    --MSA(0) <= '0' when (q1 = '1' and q0 = '0')
    --    else '1';
    with (q1 & q0 & IR) select
        MSB(1) <= '0' when "01001",
        '1' when others;
    MSB(0) <= (not q1) and Q0 and (not IR(2)) and (not IR(1)) and IR(0);
    with (q1 & q0 & IR) select
        MSC(2) <= '1' when "01010" | "01100" | "01101",
        '0' when others;
    with (q1 & q0 & IR) select
        MSC(1) <= '1' when "01100" | "01101",
        '0' when others;
```

Figure 6: VHDL for Second RALU Controller 1/2 (Part 2)

```
Text Editor - C:/QuartusProjects/Lab6/Lab6_Part2/Lab6_Part2 - [controller2.vhd]*  
File Edit View Project Processing Tools Window Help  
25 state <= Q1 & Q0;  
26  
27 IRLD <= not (Q0 or Q1);  
28 PCLD <= q1 and q0;  
29  
30 --pos 00 and 11  
31 PCINC <= (Q1 or Q0) and ((not q1) or (not q0));  
32  
33 with (q1 & q0 & IR) select  
34   MSA(1) <= '1' when "01010" | "01100" | "01101",  
35   '0' when others;  
36  
37  
38 with state select  
39   MSA(0) <= '0' when "10",  
40   '1' when others;  
41  
42 --MSA(0) <= '0' when (q1 = '1' and q0 = '0')  
43   -- else '1';0  
44  
45  
46 with (q1 & q0 & IR) select  
47   MSB(1) <= '0' when "01001",  
48   '1' when others;  
49  
50 MSB(0) <= (not q1) and Q0 and (not IR(2)) and (not IR(1)) and IR(0);  
51  
52  
53 with (q1 & q0 & IR) select  
54   MSC(2) <= '1' when "01010" | "01100" | "01101",  
55   '0' when others;  
56  
57 with (q1 & q0 & IR) select  
58   MSC(1) <= '1' when "01100" | "01101",  
59   '0' when others;  
60  
61 with (q1 & q0 & IR) select  
62   MSC(0) <= '1' when "01010" | "01101",  
63   '0' when others;  
64  
65 with (q1 & q0 & IR) select  
66   D1 <= '1' when "01000" | "01011",  
67   '0' when others;  
68  
69 D0 <= '1' when (q1 = '0' and q0 = '0') or  
70   (q1 = '0' and q0 = '1' and IR = "000")  
71   else '0';  
72  
73  
74 end behavior;  
75
```

Figure 7: VHDL for Second RALU Controller 2/2 (Part 2)

Lab 6 Report: Elementary CPU Design

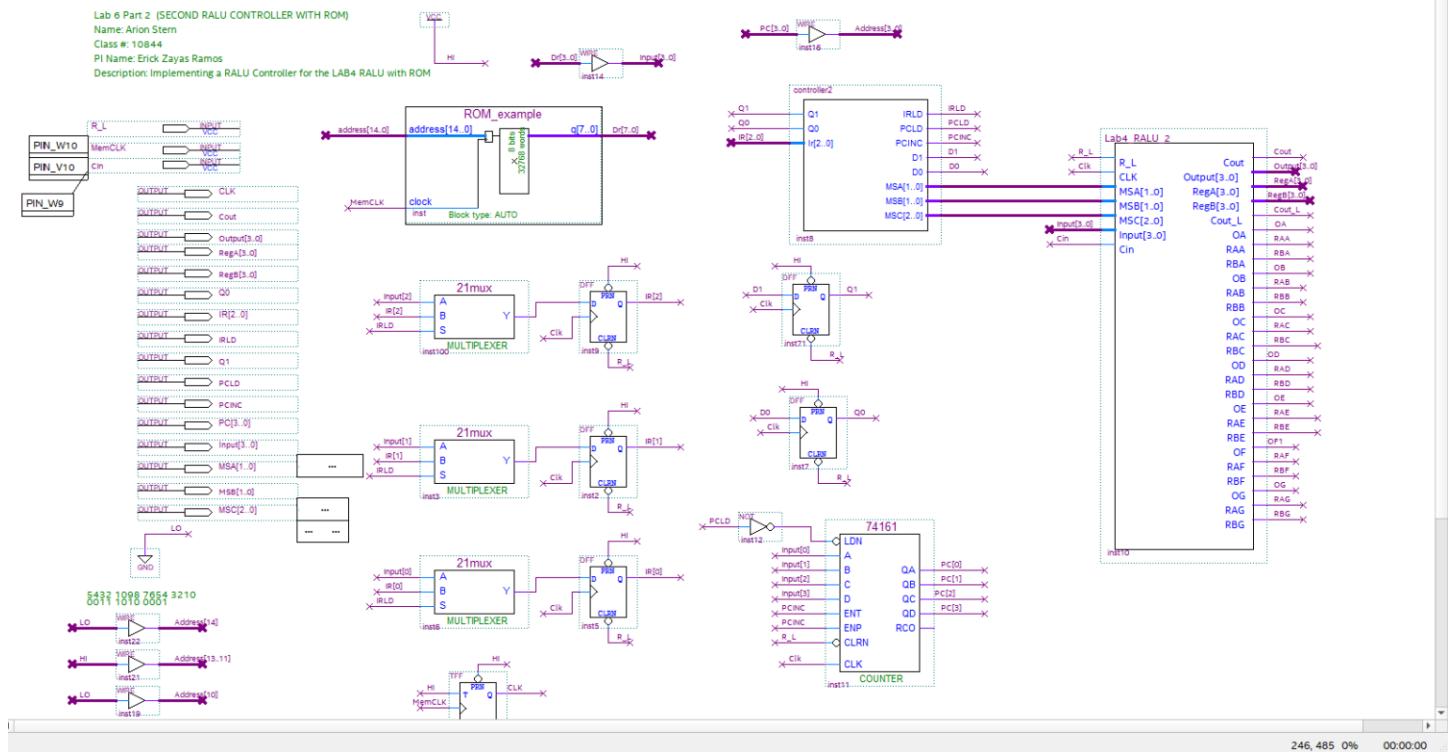


Figure 8: BDF Schematic for LAB6_Part2 1/2 (Part 2)

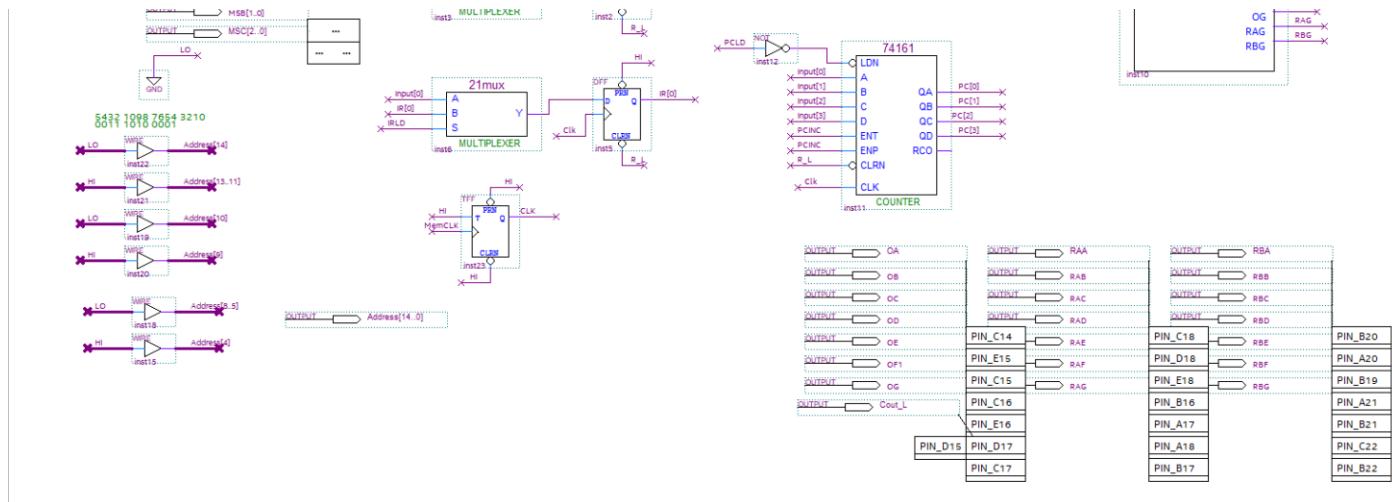


Figure 9: BDF Schematic for LAB6_Part2 2/2 (Part 2)

Addr	Instruction	Mach Codes (\$)	A	B	A	B	A	B	A	B
\$3A10	LDAA #7	3 7	7	?	X	X	X	X	X	X
\$3A12	TAB	1	7	7	X	X	X	X	X	X
\$3A13	LDAA #3	3 3	3	7	X	X	X	X	X	X
\$3A15	ABA	2	A	7	6	7	2	7	E	7
\$3A16	SAR	5	5	7	3	7	1	7	7	7
\$3A17	ABA	2	C	7	A	7	8	7	E	7
\$3A18	SAL	4	8	7	4	7	0	7	C	7
\$3A19	ABA	2	F	7	B	7	7	7	3	7
\$3A1A	JMP 5	0 5	F	7	B	7	7	7	3	7
\$3A1C	LDAA #\$F	3 F	X	X	X	X	X	X	X	X
\$3A1E	ABA	2	X	X	X	X	X	X	X	X

Table 3: Program to Assemble (Part 2)

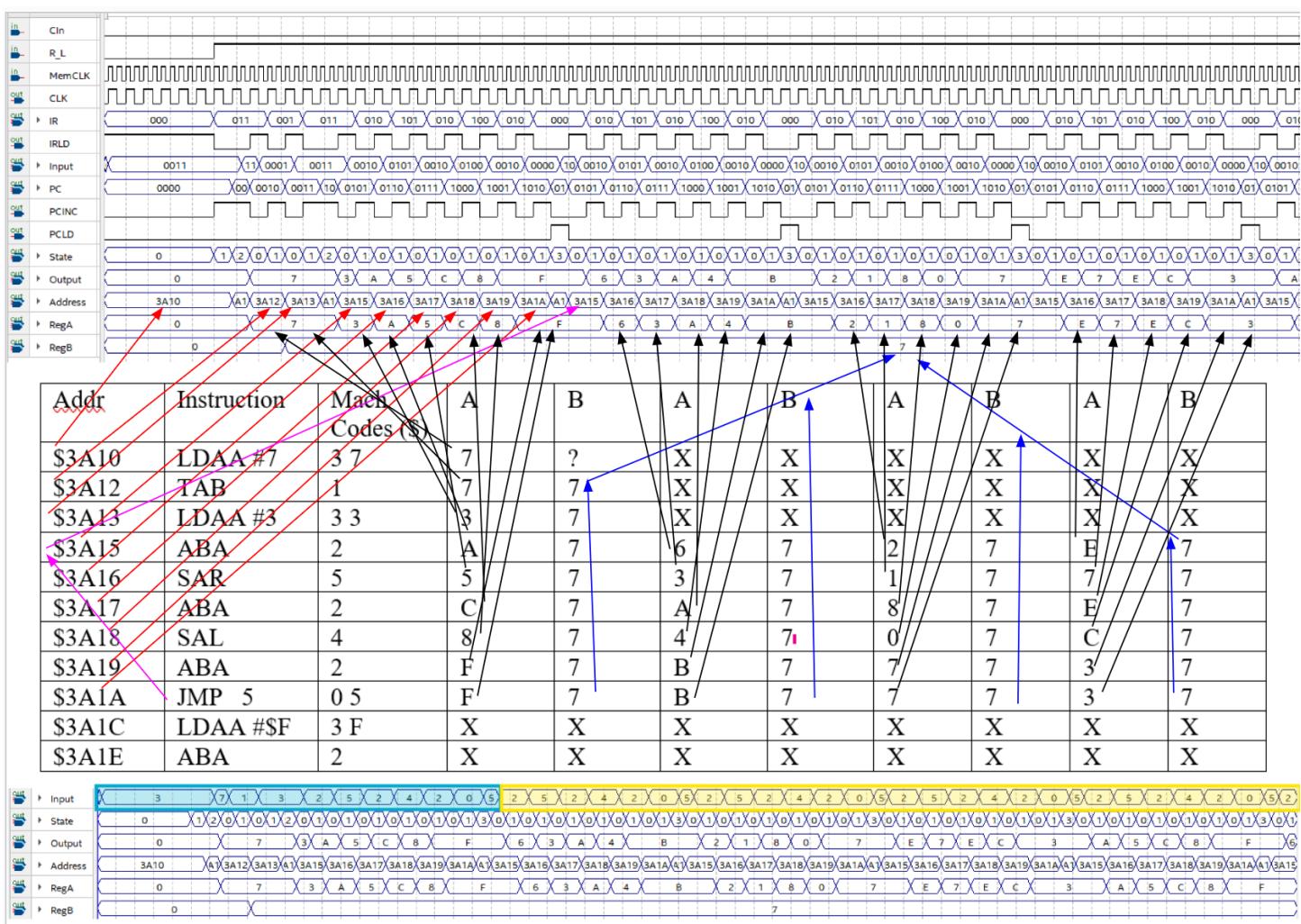
```

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13 -- refer to the applicable agreement for further details, at
14 -- https://fpgasoftware.intel.com/eula.
15
16 -- Quartus Prime generated Memory Initialization File (.mif)
17
18 WIDTH=8;
19 DEPTH=32768;
20
21 ADDRESS_RADIX=HEX;
22 DATA_RADIX=HEX;
23
24 CONTENT BEGIN
25   [0000..3A0F] : 00;
26   3A10 : 03;
27   3A11 : 07;
28   3A12 : 01;
29   [3A13..3A14] : 03;
30   3A15 : 02;
31   3A16 : 05;
32   3A17 : 02;
33   3A18 : 04;
34   3A19 : 02;
35   3A1A : 00;
36   3A1B : 05;
37   3A1C : 03;
38   3A1D : 0F;
39   3A1E : 02;
40   [3A1F..7FFF] : 00;
41 END;
42

```

Figure 10: MIF File for LAB6_Part2 ROM (Part 2)

Lab 6 Report: Elementary CPU Design



As shown above, the simulation perfectly matches my table. Also the inputs align with the machine code pattern of 3 7 1 3 3 2 5 2 4 2 0 5, with the highlighted sequence of numbers repeating due to the jump.

Figure 11: Annotated Simulation for LAB6_Part2 (Part 2)