

Data Structures & Algorithms

Adil M. Khan
Professor of Computer Science
Innopolis University
Kazan, Russia

Trees

A Question!

- Last time, we learned that a hash table can insert and retrieve elements in $O(1)$
- This is as fast as it can be!
- Why do we need Trees then?
- Stay patient, we will discuss this after we are done with Binary Search Trees

Topic Overview

- Trees
- Oriented Trees
- Ordered Trees
- Binary Trees
- Tree Traversal Algorithms
- Implementation

Tree

- A tree combines the advantages of two other data structures:
 - An ordered array
 - A linked list

Ordered Array

- Quick to search for a particular element, using binary search
- On the other hand, insertions are slow
 - First need to find where the object will go
 - Then move all the objects with greater keys up one space to make the room

Linked List

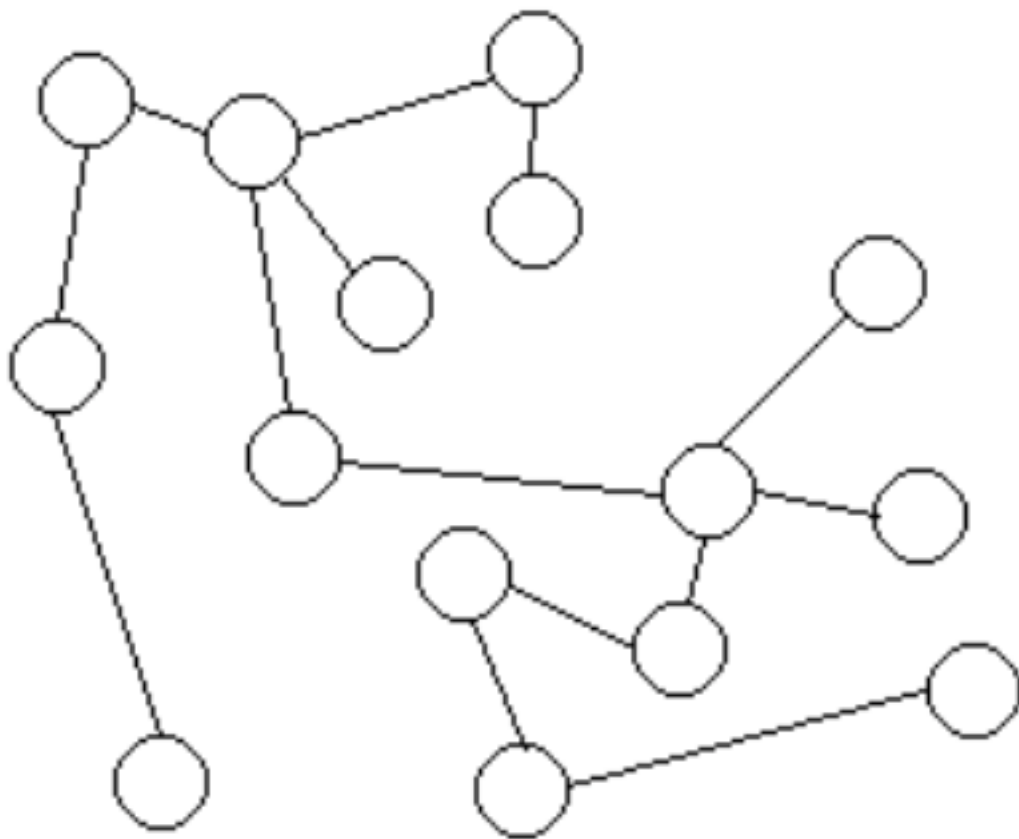
- Insertions and deletions are quick
 - Simply requires changing a few constant number of references
- On the other hand, finding a particular element is slow
 - Must start at the beginning of the linked list
 - Visit each element until you find the one you're looking for.
- What if we made the linked list ordered? Will it help?

Trees to Rescue

- It would be nice to have a data structure
 - With quick insertions and deletions of a linked list
 - And, the quick searching of an ordered array
- Trees provide both these characteristics
- Our main focus will be a **binary tree**
- But let's first start discussing trees in general

Tree

- Consists of nodes connected by edges



Node

(a.k.a. vertex) a data element in the tree

Edge

(a.k.a. branch, or link) a connection between two nodes

Empty tree

has zero nodes

Size

the size of a tree is the number of nodes

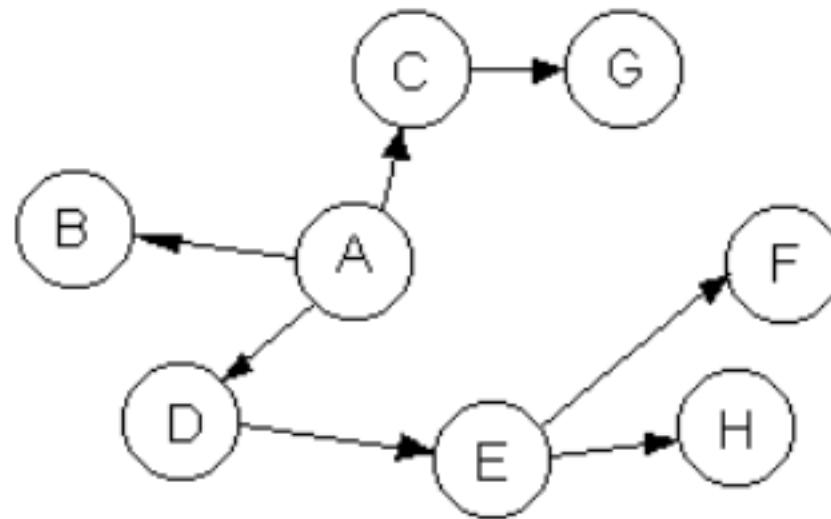
Path

a sequence $n_0e_1n_1e_2n_2...e_kn_k$ where $k \geq 0$ and e_i connects n_{i-1} and n_i .

Oriented Trees

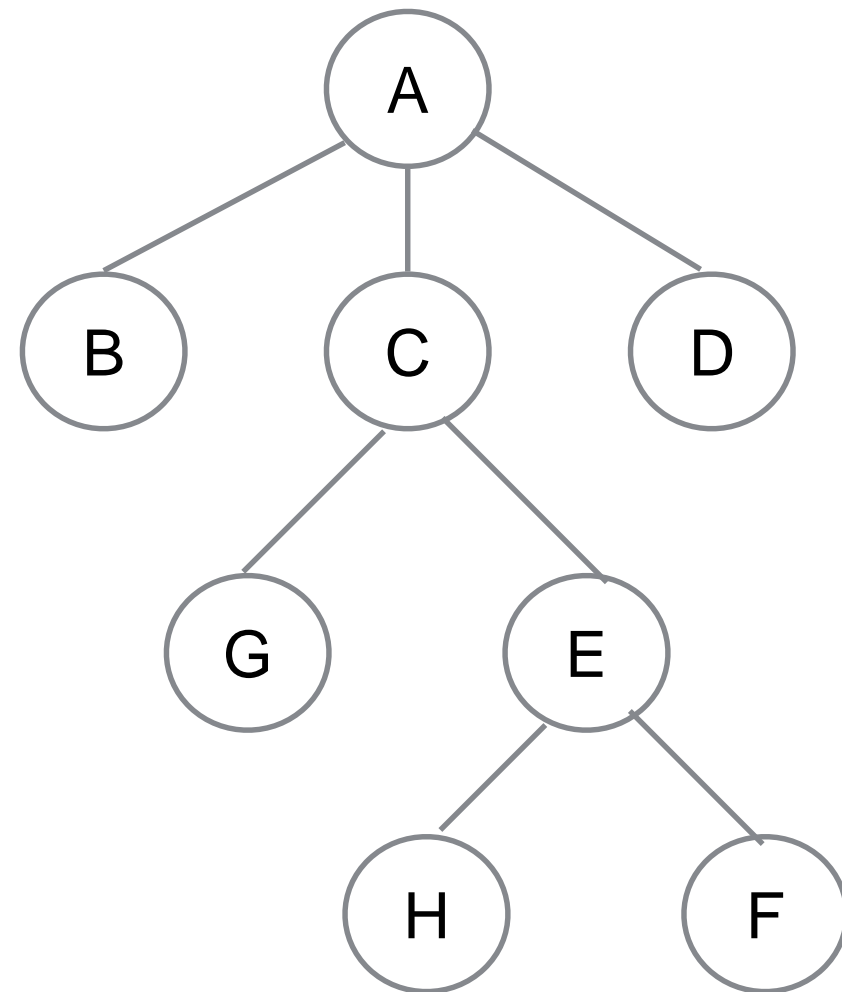
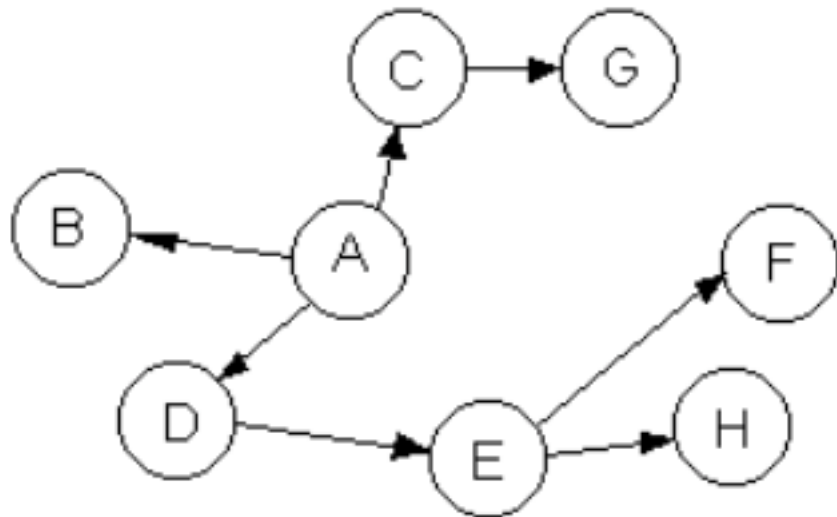
Oriented Trees

- A tree used to represent a hierarchical data.
- All edges are directed outward from a distinguished node called the root node



Oriented Trees

- Usually drawn with the root at the top, all edges pointing downward, the arrows are thus redundant and often omitted.



Definitions

Parent

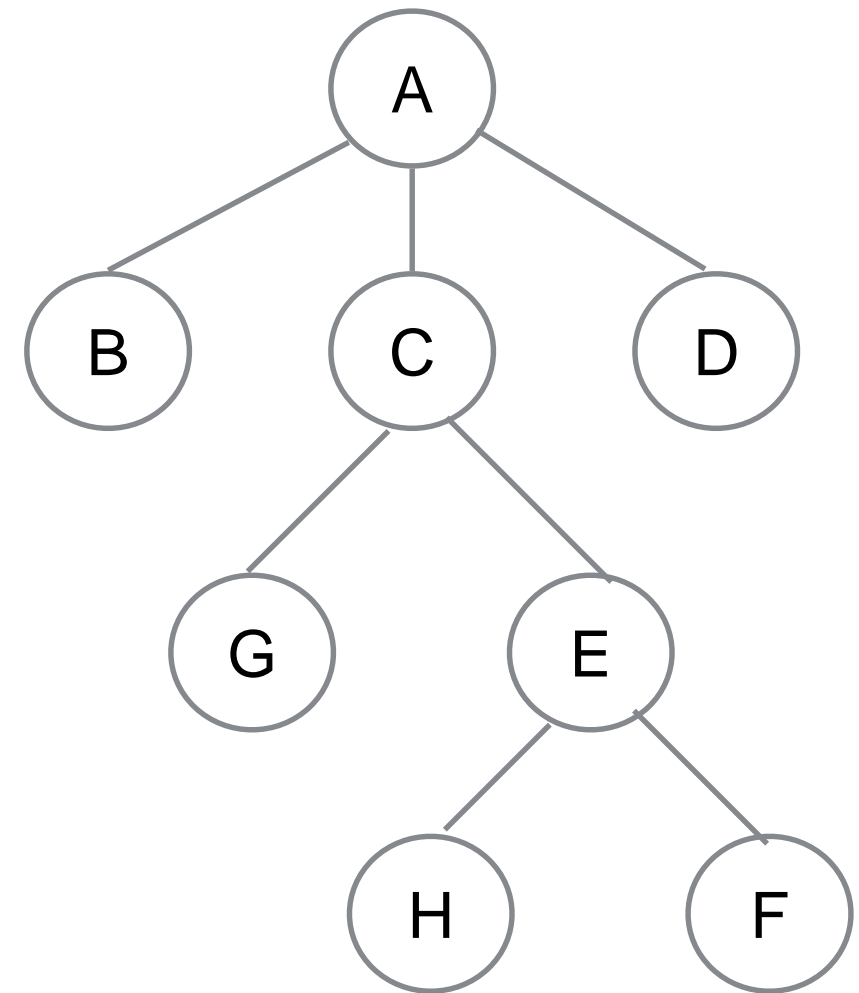
A is the parent of B and C and D

Children

The children of A are B and C and D

Siblings

B and C are siblings



Definitions

Root node

The only node without a parent

Leaf node

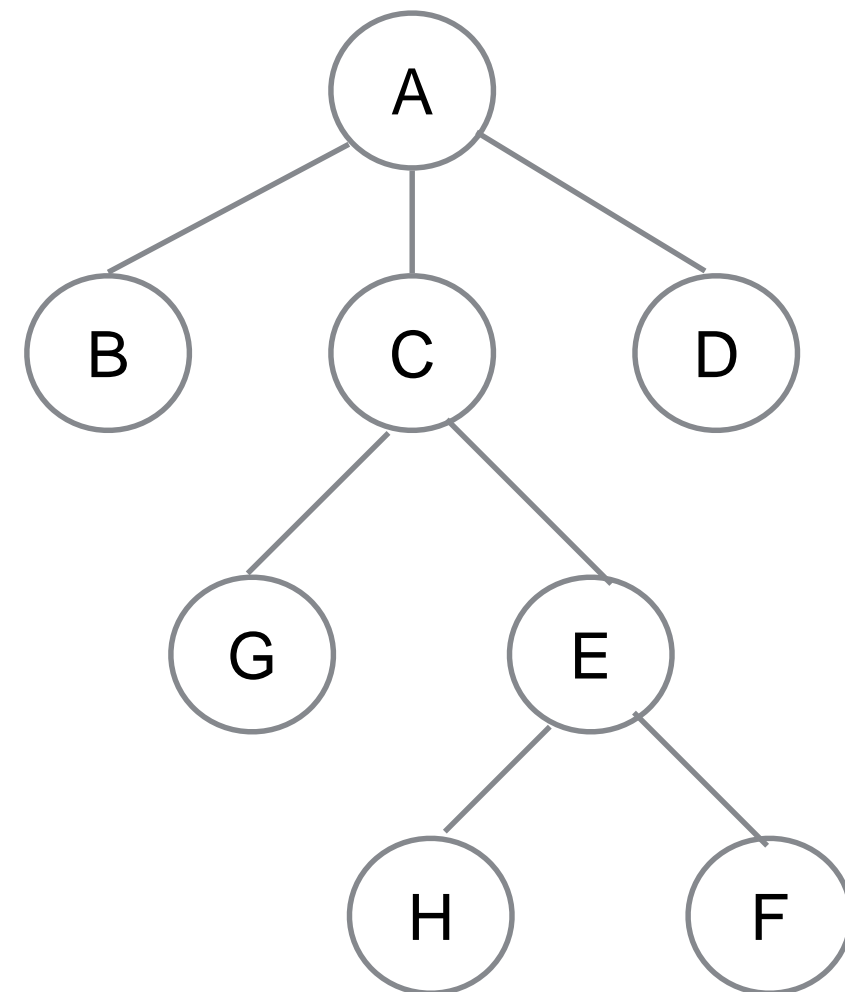
A node without children

Degree(n)

The number of children of n

Degree(t)

The greatest degree of the nodes in t



Definitions

Level(n)

$\text{Level}(n) = \text{if } n \text{ is the root then } 0 \text{ else } 1 + \text{Level}(\text{Parent}(n))$

Depth(n)

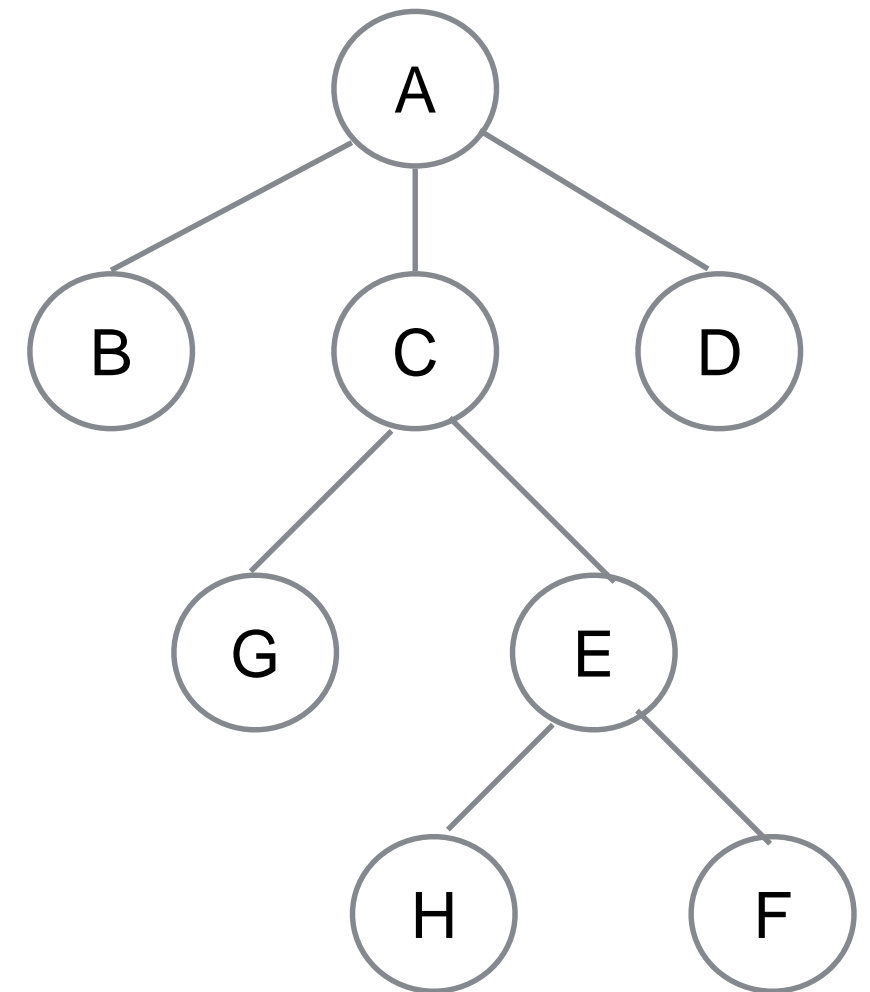
Same as $\text{Level}(n)$

Height(n)

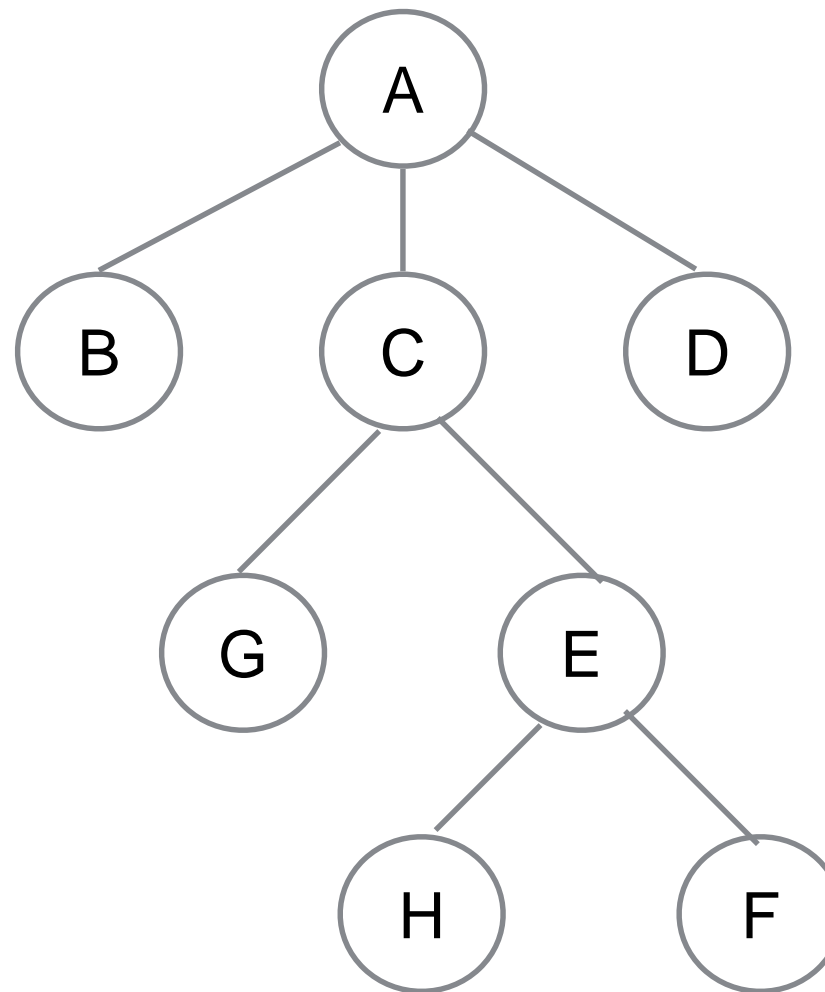
Maximum length over all paths from n to a leaf

Height(t)

The height of t 's root node



Definitions



Ancestors(n)

$\text{Ancestors}(n) = \text{if } n \text{ is the root then } \{ \} \text{ else } \{ \text{Parent}(n) \} \text{ union } \text{Ancestors}(\text{Parent}(n))$

Descendants(n)

The set of all nodes reachable from n following the edges leaving it in the direction of the arrows

Definitions

PathLength(t)

The sum of all the depths of all the nodes in t

Empty Tree

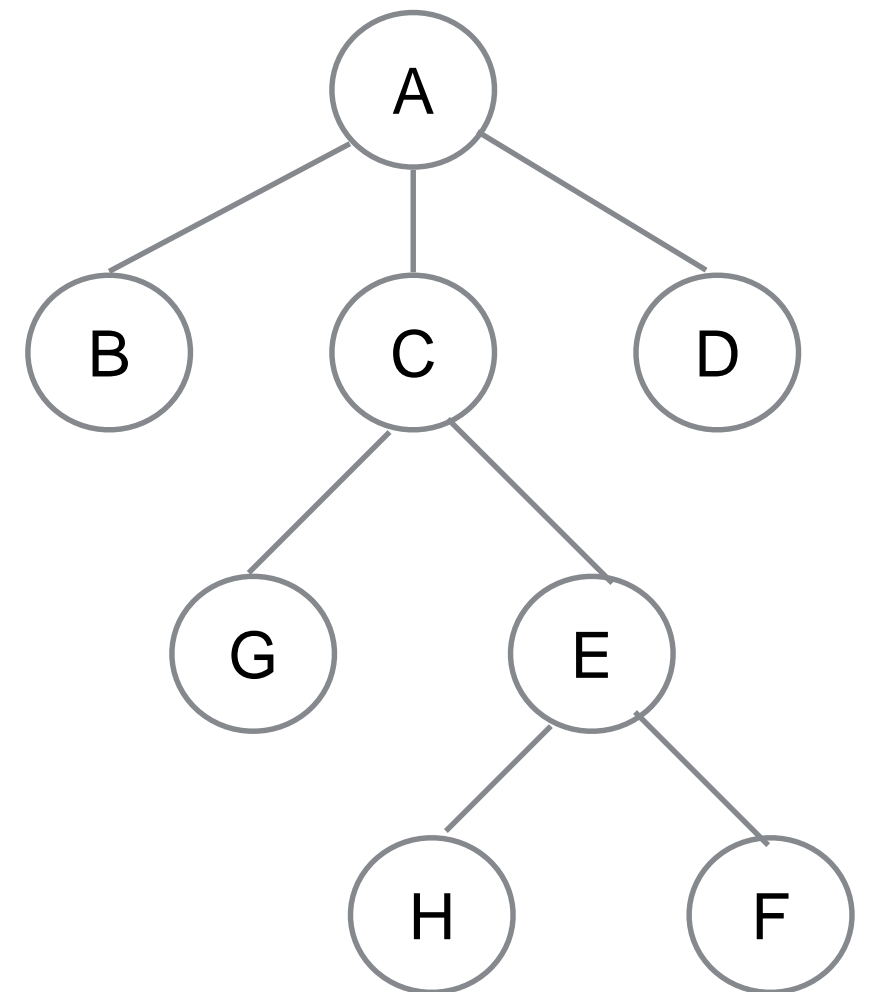
A tree with zero nodes

Singleton Tree

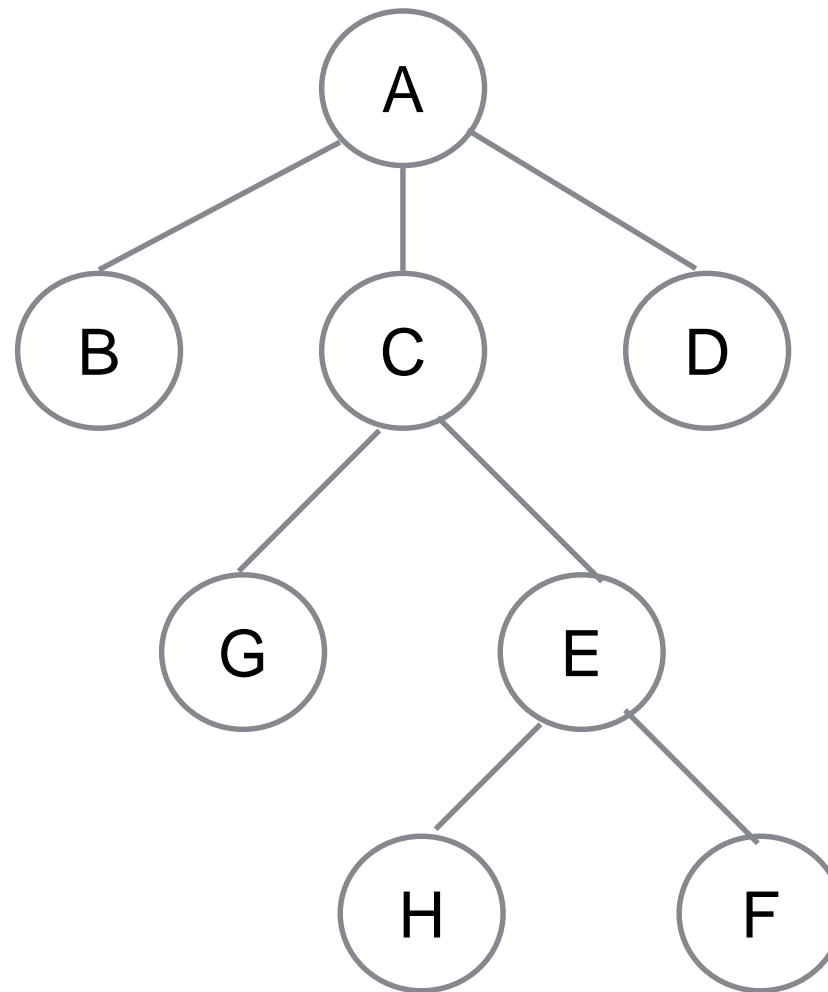
A tree with only one node

Subtree(n)

The subtree rooted at n



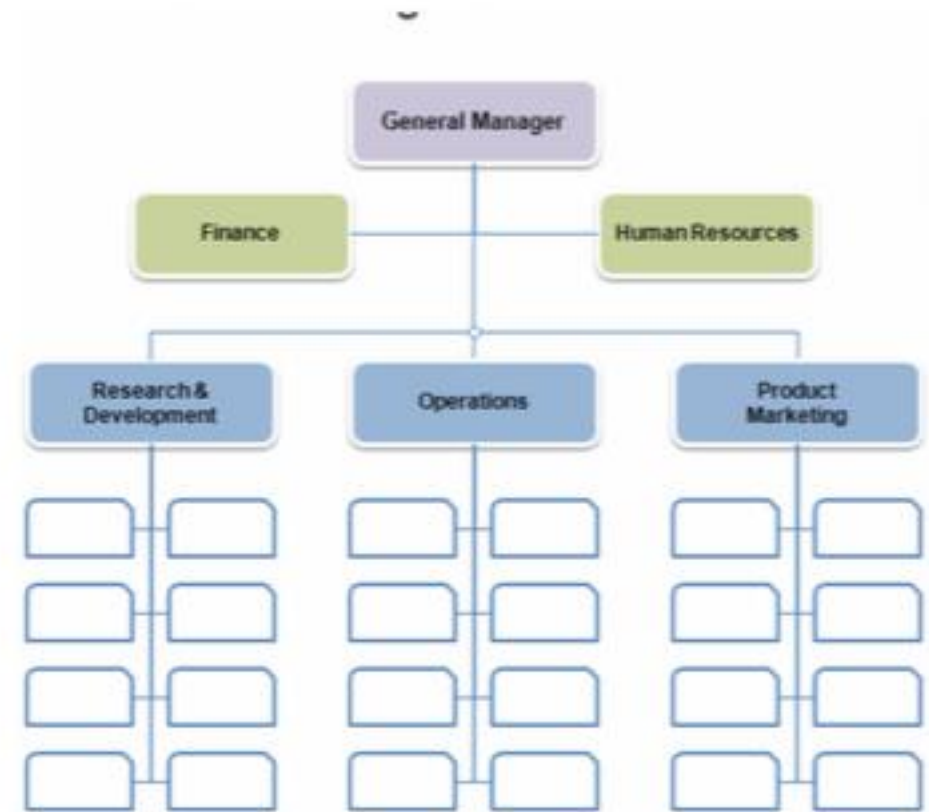
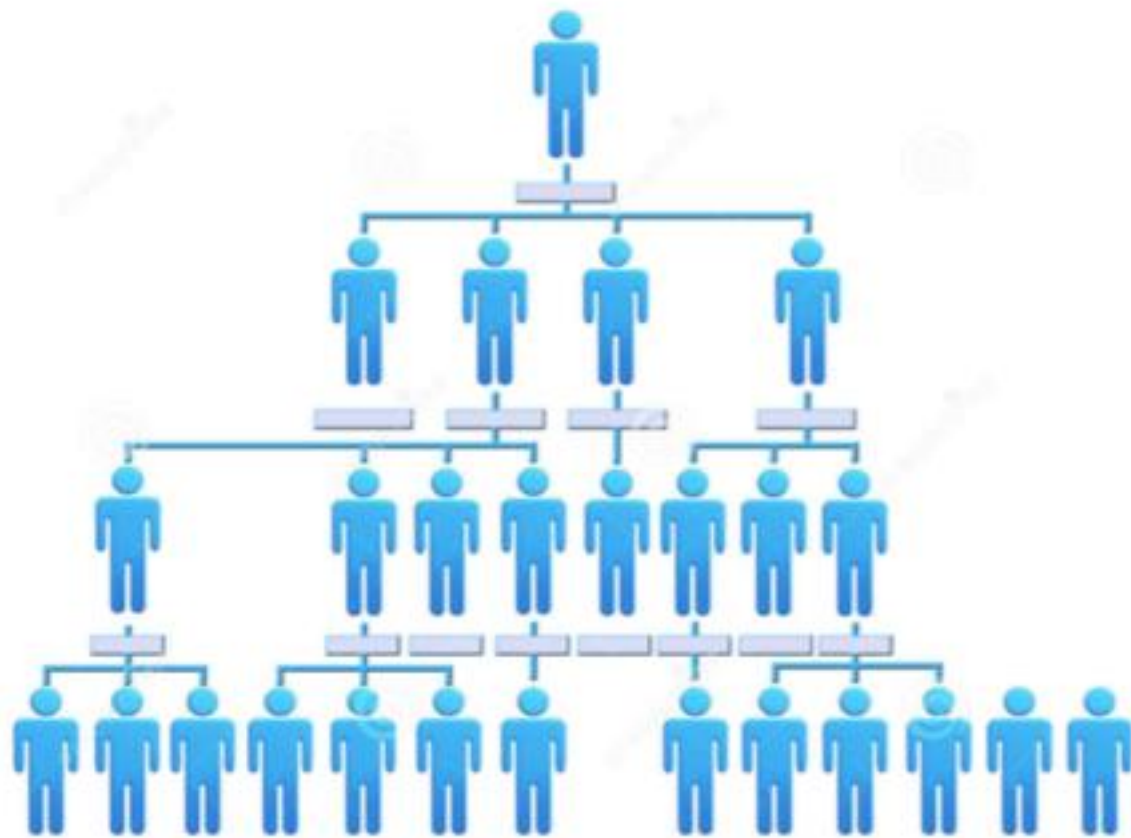
Definitions



k-node

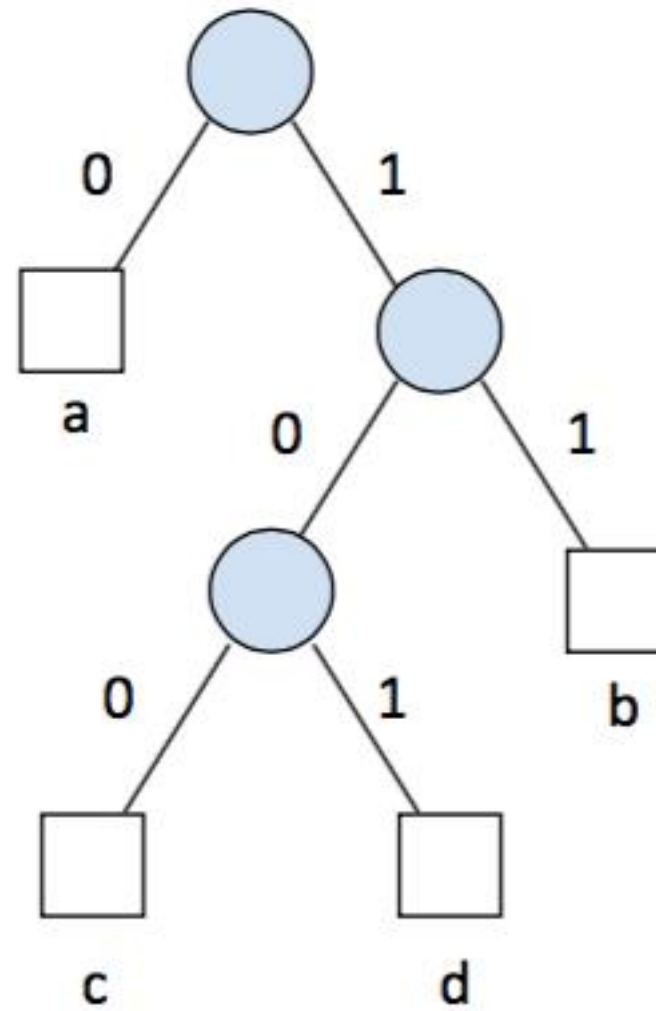
A node with k children. A 0-node is a leaf. A 1-node has exactly one child. A 2-node has exactly two children. Etc.

Uses of Trees



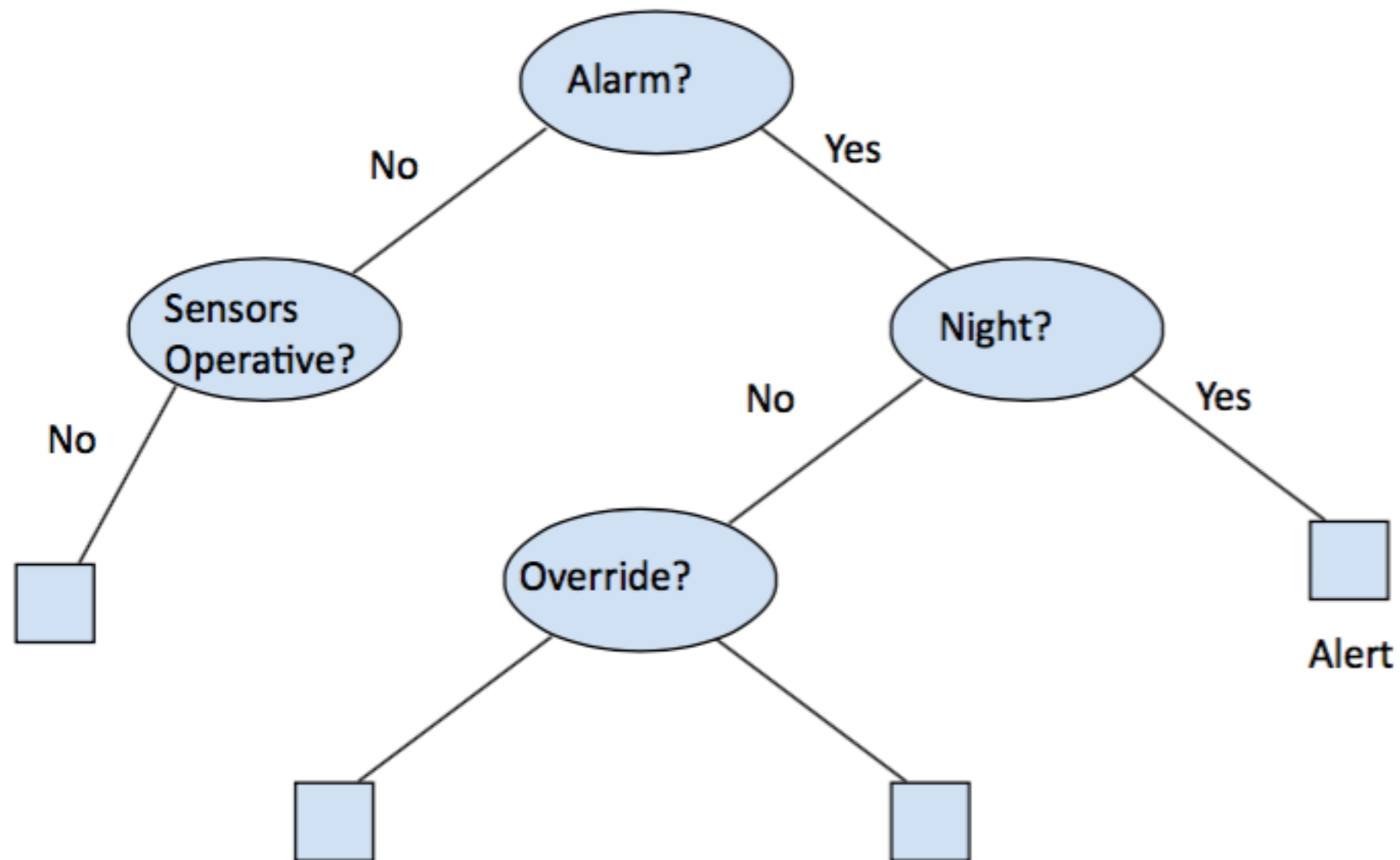
Organization Tree

Uses of Trees



Code Tree

Uses of Trees

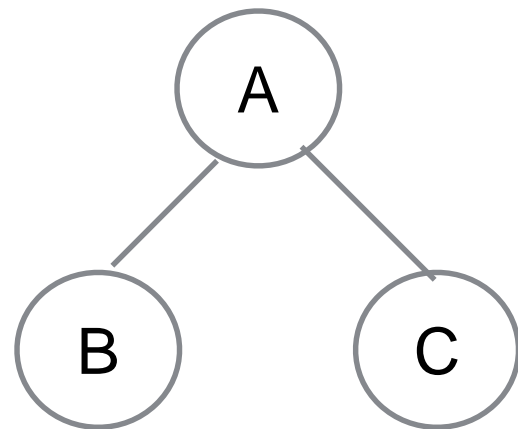


Decision Trees

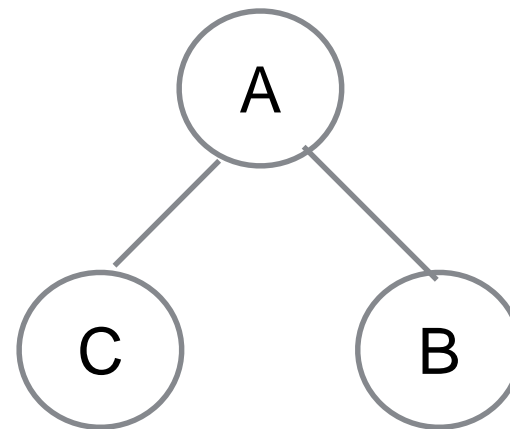
Ordered Trees

Ordered Trees

- An oriented tree in which the children of a node are somehow **ordered**.



T1



T2

If T1 and T2 are ordered trees then $T1 \neq T2$,
otherwise $T1 = T2$

Types of Ordered Trees

- Fibonacci Tree
- Binomial tree
- k-ary Tree

Fibonacci Trees

- A Fibonacci Tree (F_k) is defined by
 - F_0 is the empty tree
 - F_1 is a tree with only one node
 - F_{k+2} is a node whose left subtree is a F_{k+1} tree and whose right subtree is a F_k tree.

Exercise: Draw the trees F_0 through F_5 .

Binomial Trees

- The Binomial Tree (B_k) consists of a node with k children. The first child is the root of a B_{k-1} tree, the second is the root of a B_{k-2} tree, etc.

Exercise: Draw the trees B_0 through B_5 .

k-ary Trees

- A tree in which the children of a node appear at distinct index positions in $0..k-1$
- Maximum number of children for a node is k

Types of k-ary Trees

- 2-ary Trees, known as **Binary Trees**
- 3-ary Trees, known as **Ternary Trees**
- 1-ary Trees, known as **Lists**

k-ary Tree Representation

```
class Node {
```

```
    private Object data;
```

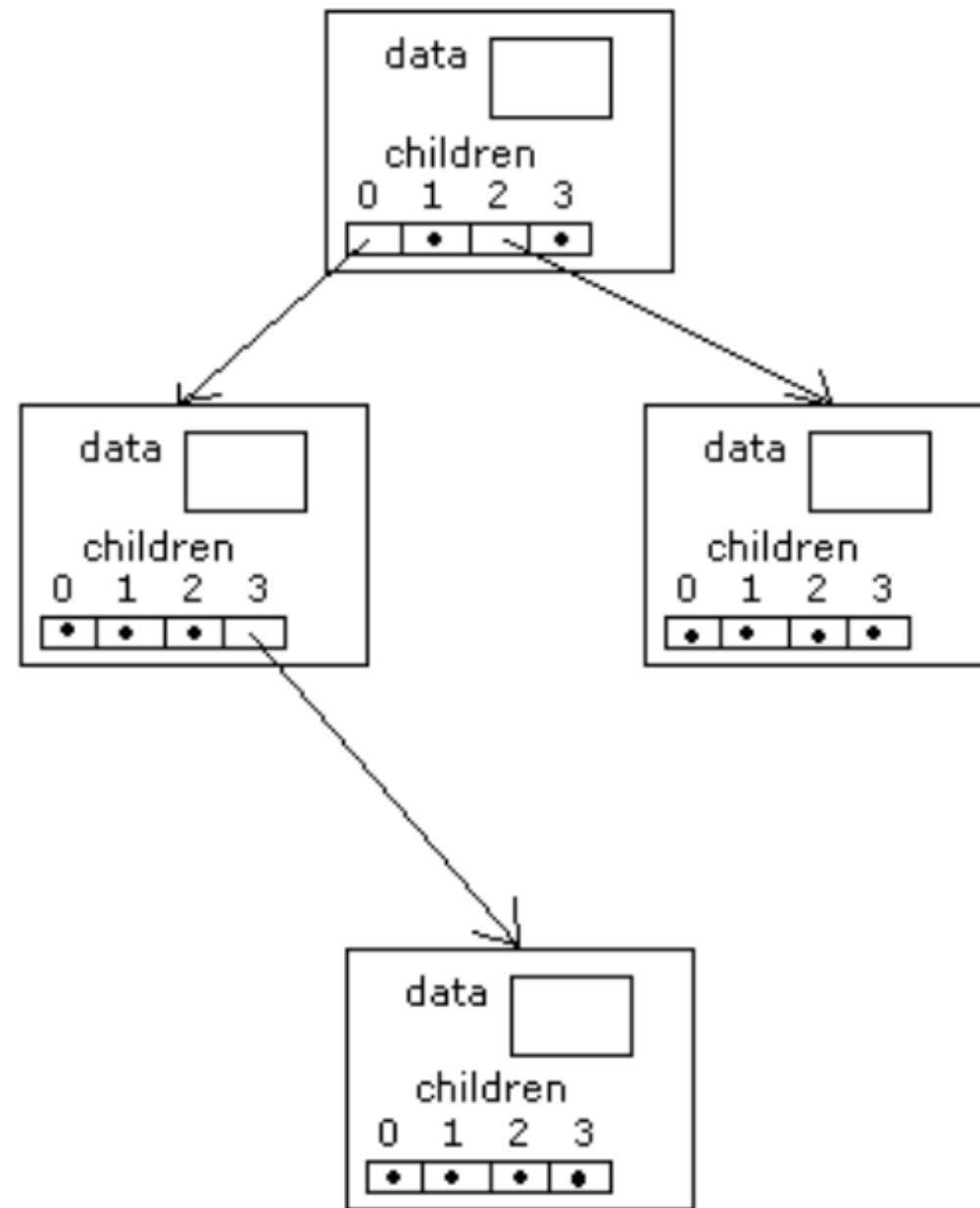
```
    private Node[] children;
```

```
    .
```

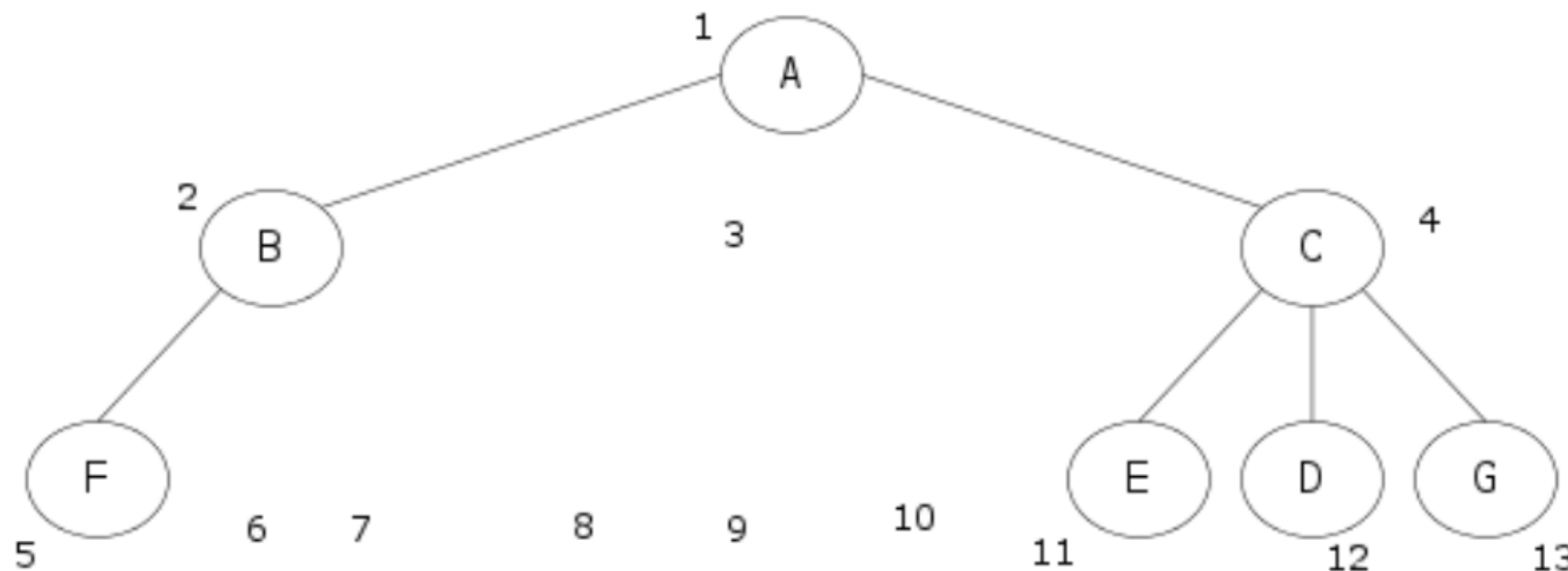
```
    .
```

```
    .
```

```
}
```



Array Representation



1	A
2	B
3	
4	C
5	F
6	
7	
8	
9	
10	
11	E
12	D
13	G

In a binary tree

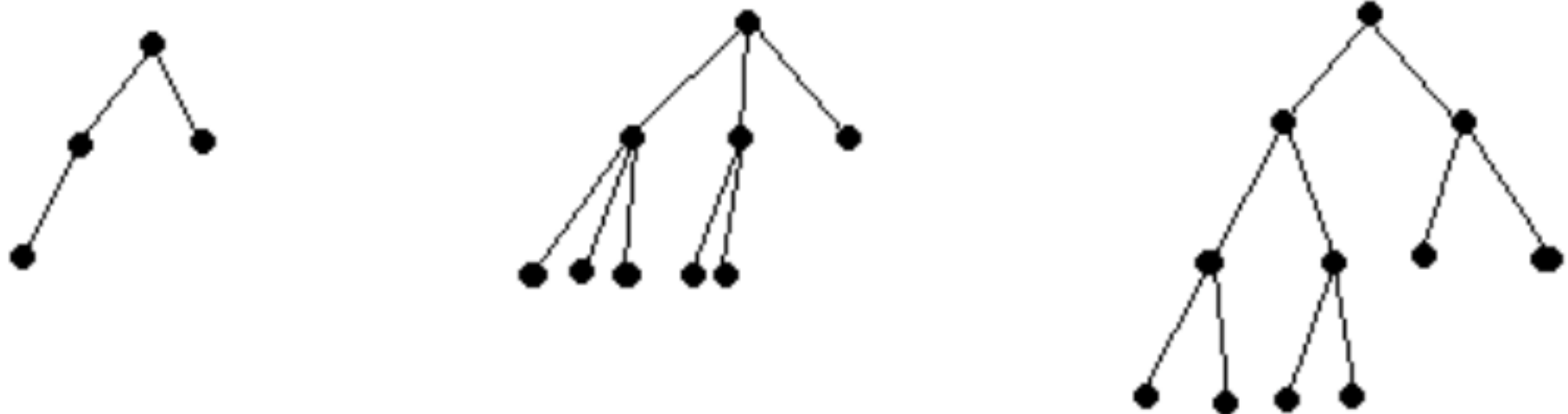
- $\text{parent}(i)$ is $i / 2$
- $\text{leftChild}(i)$ is $2i$
- $\text{rightChild}(i)$ is $2i + 1$

In a k-ary tree

- $\text{parent}(i)$ is $(k + i - 2) / k$
- $j\text{th child of } i$ is $ki - (k-2) + j$

Complete k-ary Trees

- Filled out on every level, except perhaps on the last one
- All nodes on the last level, should be as far to left as possible



Complete trees can be packed into an array with no wasted space

Binary Tree

Binary Tree

- A **binary tree** T of n nodes, $n \geq 0$,
 - either is empty, if $n = 0$
 - or consists of a **root node** u and two binary trees $u(1)$ and $u(2)$ of n_1 and n_2 nodes, respectively, such that $n = 1 + n_1 + n_2$
- We say that $u(1)$ is the **first or left subtree** of T , and $u(2)$ is the **second or right subtree** of T

Binary Tree



External nodes - have no subtrees (referred to as leaf nodes)



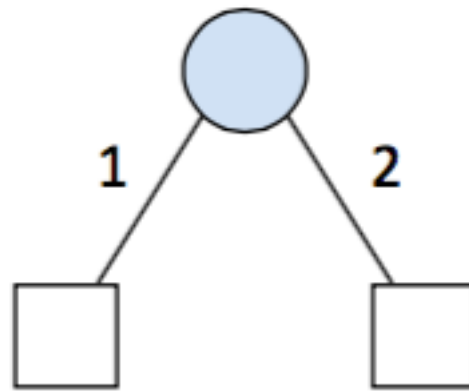
Internal nodes - always have two subtrees

Binary Tree



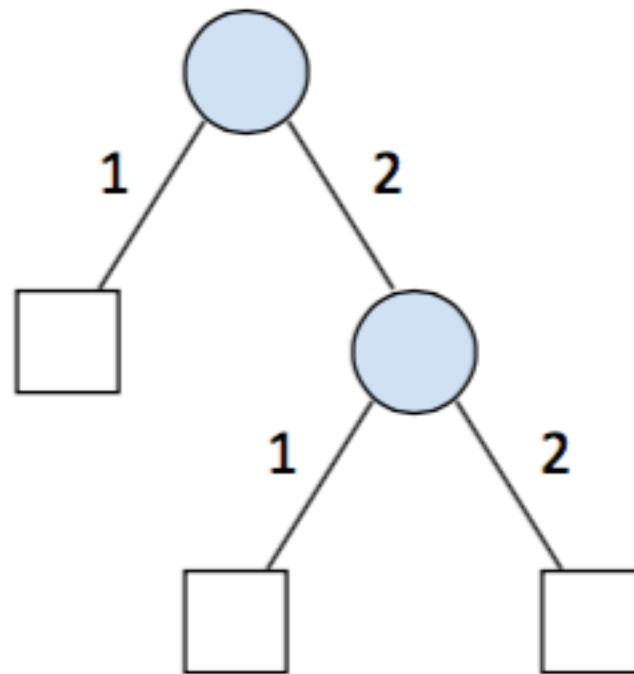
Binary Tree of zero nodes

Binary Tree



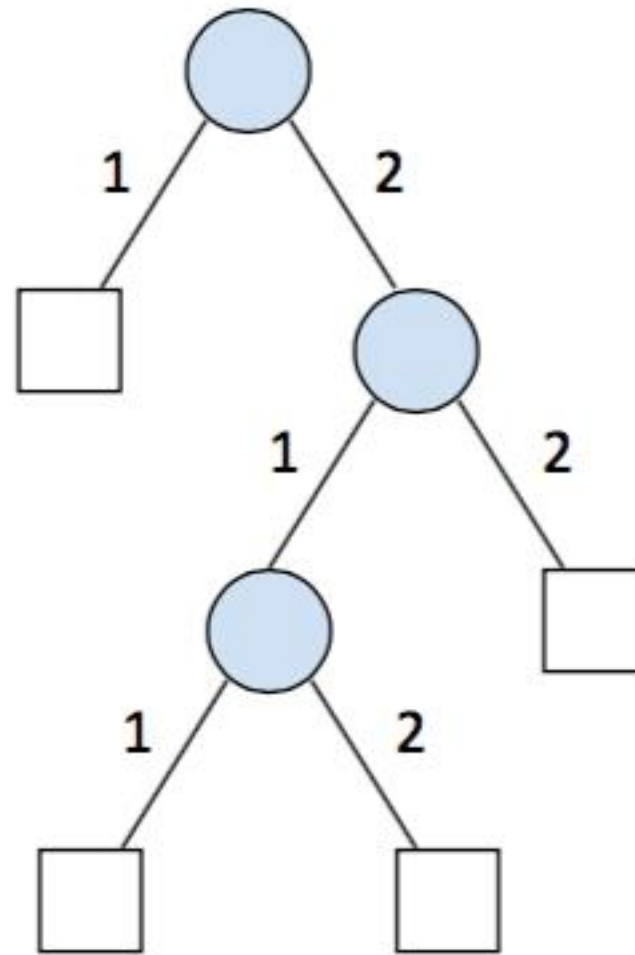
Binary Tree of 1 nodes

Binary Tree



Binary Tree of 2 nodes

Binary Tree



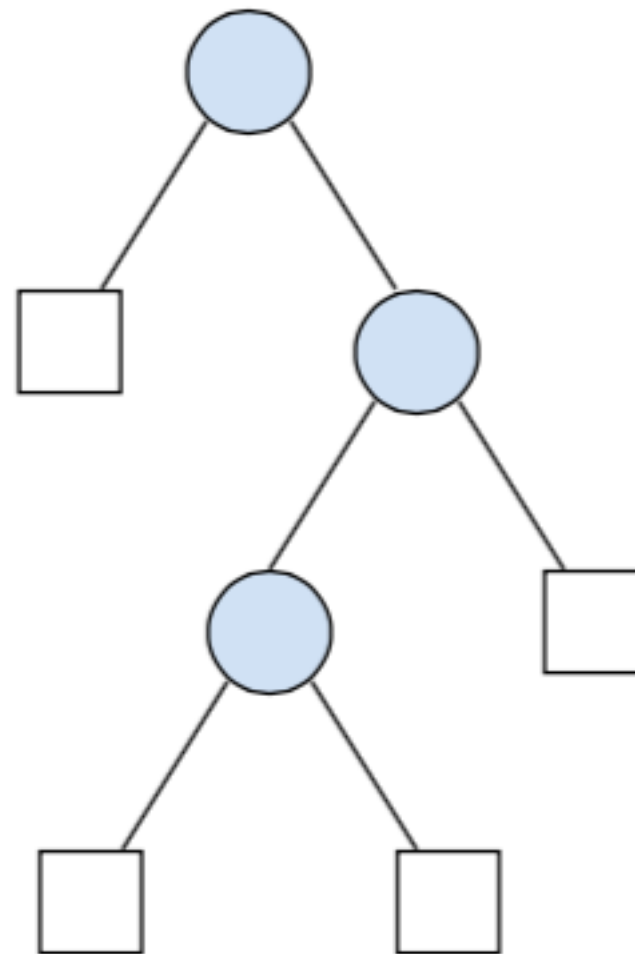
Binary Tree of 3 nodes

Why Binary Trees

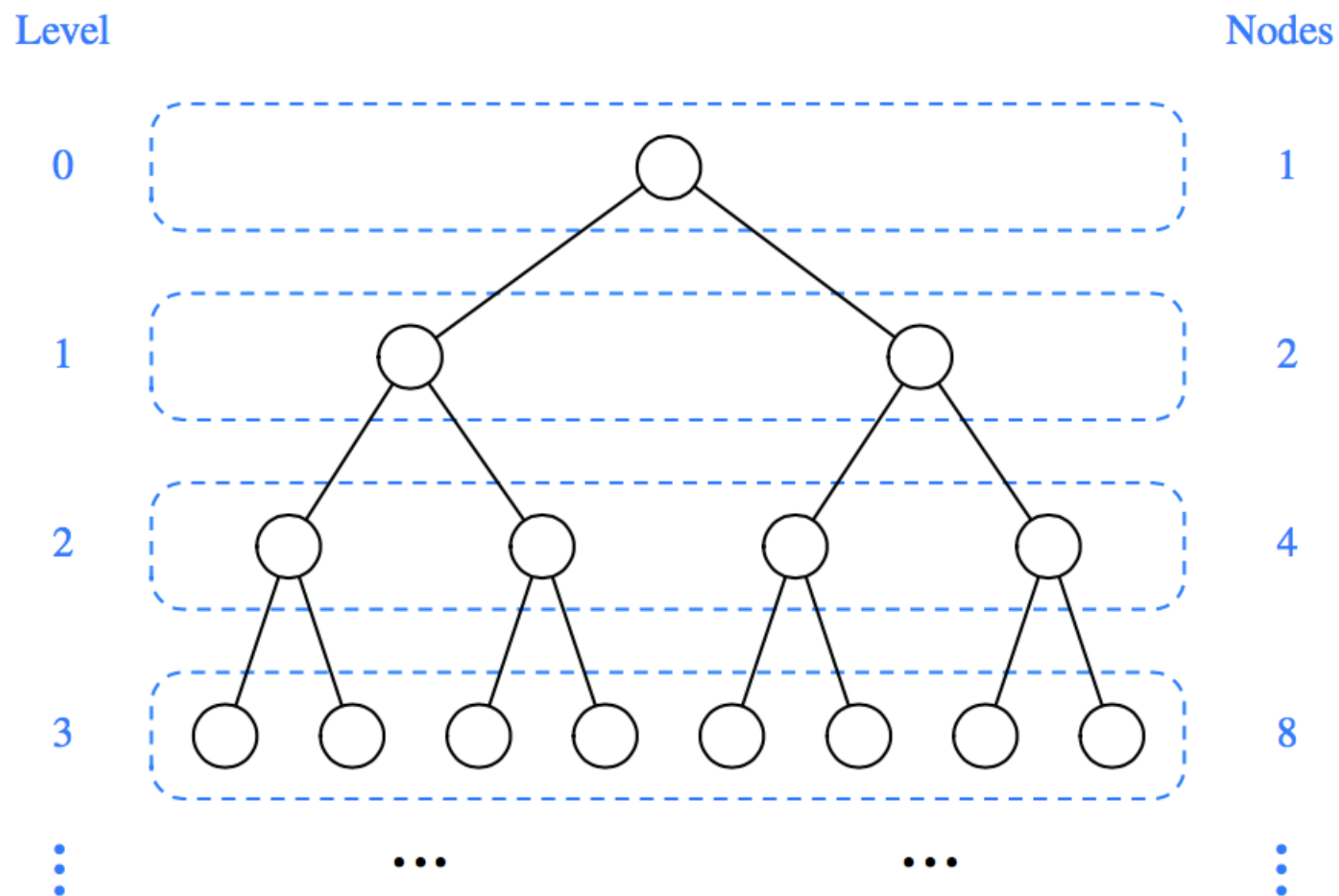
- Binary trees are a bit simpler and easier to understand than trees with a large or unbounded number of children
- Binary tree nodes have an elegant linked representation with "left" and "right" subtrees
- Binary trees form the basis for efficient representations of sets, dictionaries, and priority queues

Binary Tree Properties

- Let T be a binary tree of size n , $n \geq 0$,
- Then, the number of external nodes of T is $n+1$



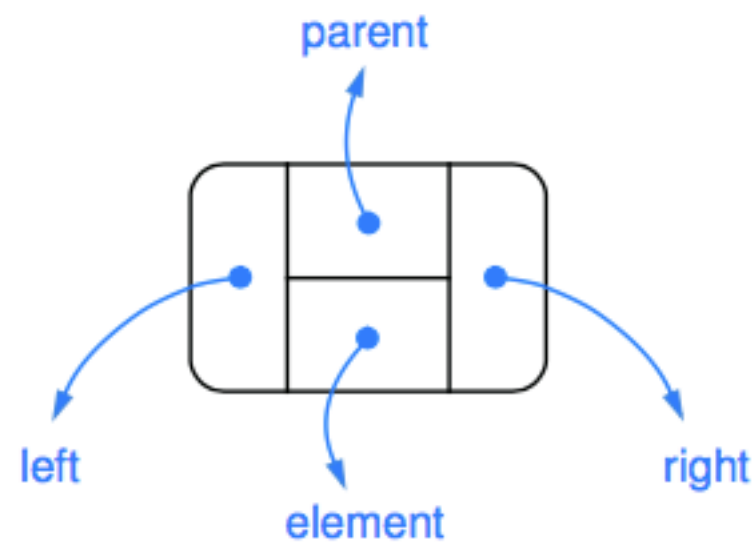
Binary Tree Properties



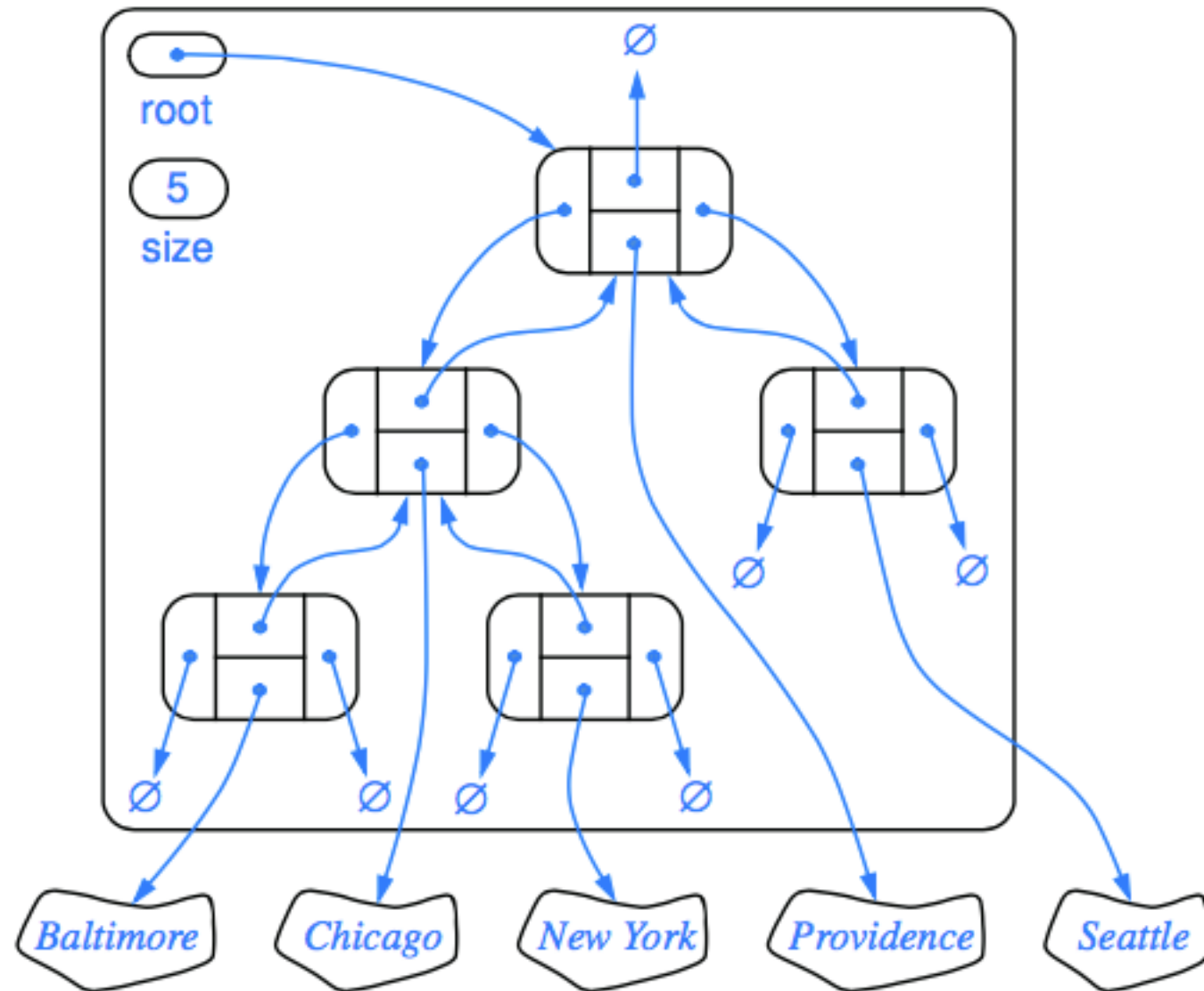
Maximum number of nodes in the levels of a binary tree

Binary Tree Representations

- Linked Structure



(a)



(b)

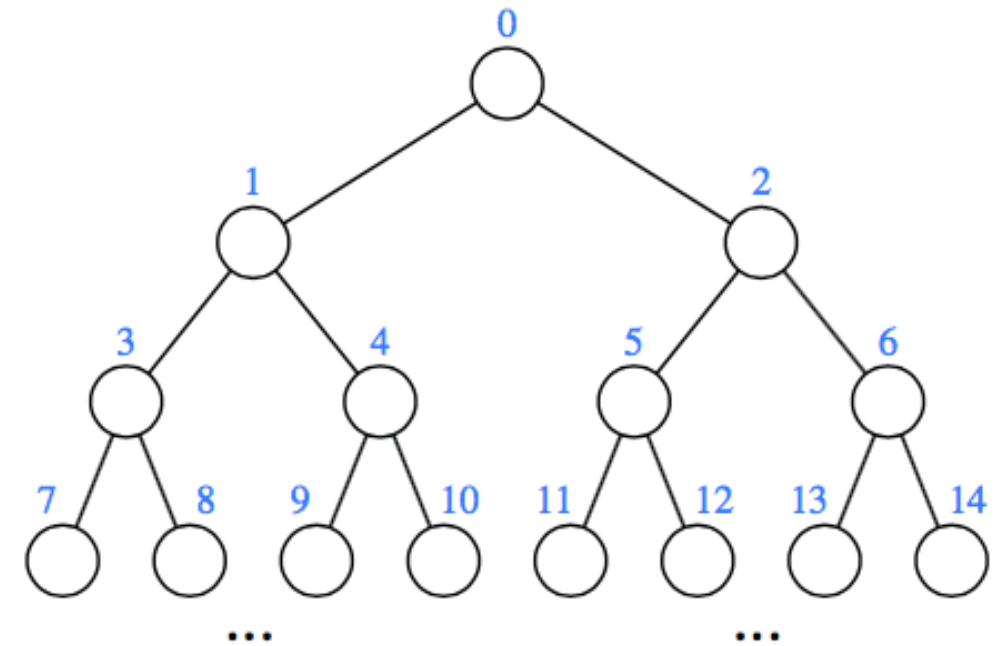
Binary Tree Representations

- Array Based Structure
 - Requires a mechanism for numbering the positions of T
 - For every position p of T , let $f(p)$ be the integer defined as follows
 - If p is the root of T , then $f(p) = 0$
 - If p is the left child of position q , then $f(p) = 2f(q)+1$.
 - If p is the right child of position q , then $f(p) = 2f(q)+2$.

Binary Tree Representations

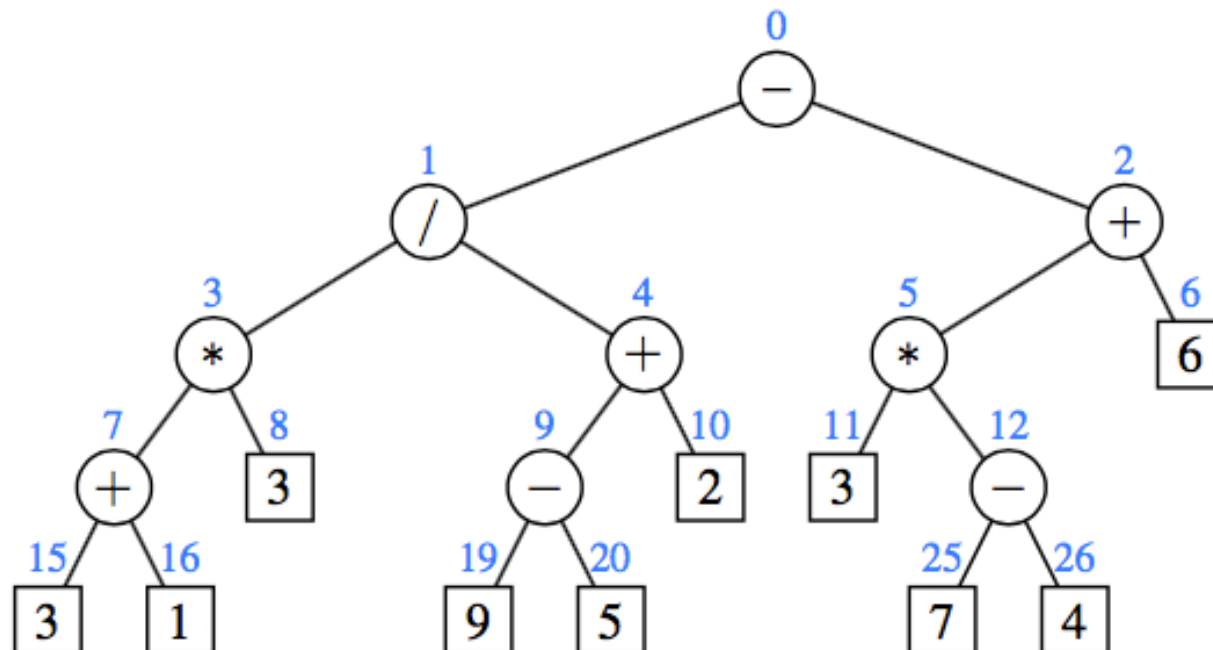
- Array Based Structure

(a)



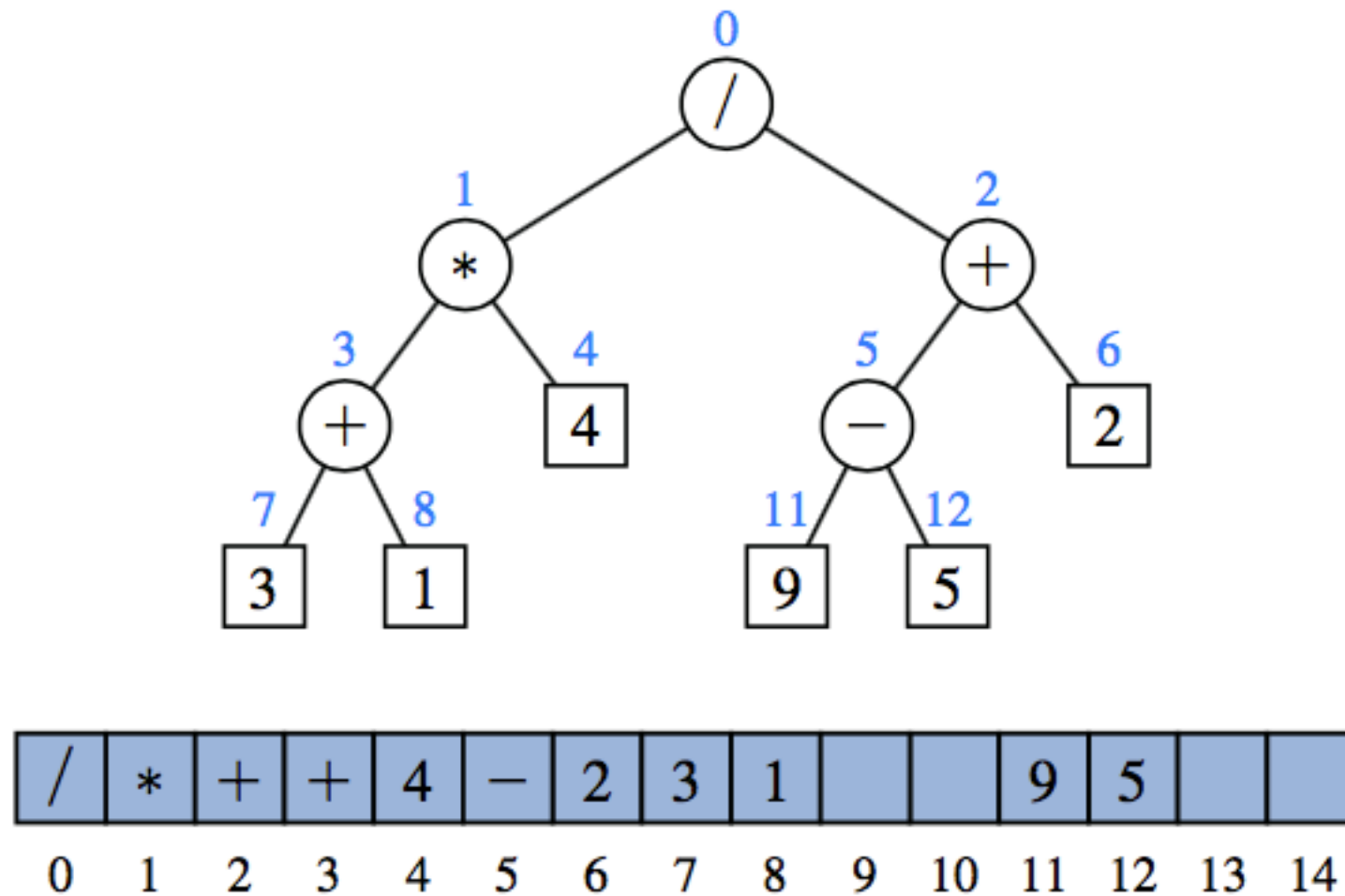
(a) general scheme; (b) an example.

(b)



Binary Tree Representations

- Array Based Structure

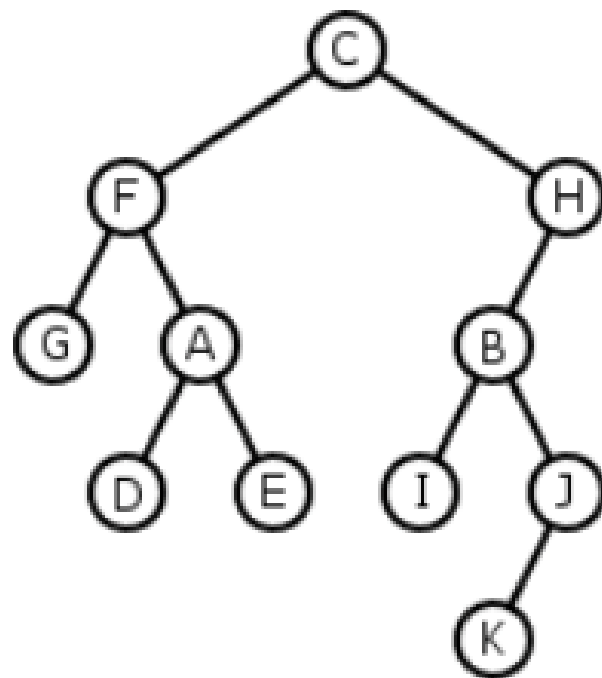


Tree Traversal Algorithms

Tree Traversal Algorithms

- A way of **accessing** or **visiting** all the nodes of T
- Preorder Traversal –
 - Visit the node, Preorder Left, Preorder Right
- Postorder Traversal
 - Postorder Left, Postorder Right, Visit the node
- Inorder Traversal
 - Inorder Left, Visit the node, Inorder Right

Tree Traversals



- **Preorder:** C F G A D E H B I J K
- **Inorder:** G F D A E C I B K J H
- **Postorder:** G D E A F I K J B H C

Preorder Traversal

Algorithm preorder(p):

perform the “visit” action for position p { this happens before any recursion }

for each child c in $\text{children}(p)$ **do**

```
preorder(c) { recursively traverse the subtree rooted at c }
```

Code Fragment 8.12: Algorithm preorder for performing the preorder traversal of a subtree rooted at position p of a tree.

Postorder Traversal

Algorithm postorder(p):

for each child c in $\text{children}(p)$ do

```
postorder(c) { recursively traverse the subtree rooted at c }
```

perform the “visit” action for position p { this happens after any recursion }

Code Fragment 8.13: Algorithm `postorder` for performing the postorder traversal of a subtree rooted at position p of a tree.

Inorder Traversal

Algorithm inorder(p):

if p has a left child lc then

$$\text{inorder}(lc)$$

{ recursively traverse the left subtree of p }

perform the “visit” action for position p

if p has a right child rc then

$$\text{inorder}(rc)$$

{ recursively traverse the right subtree of p }

Code Fragment 8.15: Algorithm inorder for performing an inorder traversal of a subtree rooted at position p of a binary tree.

Implementation

Linked Binary Tree Node

```
1 public class LinkedBinaryTreeNode<E>{
2
3     protected E data;
4     protected LinkedBinaryTreeNode<E> parent;
5     protected LinkedBinaryTreeNode<E> left;
6     protected LinkedBinaryTreeNode<E> right;
7
8     /**
9      * Constructs a node with a given element
10     */
11     public LinkedBinaryTreeNode(E data) {
12         this.data = data;
13         this.parent = null;
14         this.left = null;
15         this.right = null;
16     }
17 }
```

Linked Binary Tree Node

```
18  /**
19   * Constructs a node with a given element and neighbors
20   */
21  public LinkedBinaryTreeNode(E data, LinkedBinaryTreeNode<E> parent,
22     LinkedBinaryTreeNode<E> leftChild, LinkedBinaryTreeNode<E> rightChild) {
23     this.data = data;
24     this.parent = parent;
25     this.left = leftChild;
26     this.right = rightChild;
27 }
28
29  /**
30   * Returns the data stored in this node.
31   */
32  public E getData() {
33     return data;
34 }
35
```

Linked Binary Tree Node

```
36  /**
37   * Modifies the data stored in this node.
38   */
39  public void setData(E data) {
40      this.data = data;
41  }
42
43  /**
44   * Returns the parent of this node, or null if this node is a root.
45   */
46  public LinkedBinaryTreeNode<E> getParent() {
47      return parent;
48  }
49
50  /**
51   * Sets the parent of this node.
52   */
53  public void setParent(LinkedBinaryTreeNode<E> parent) {
54      this.parent = parent;
55  }
56
```

Linked Binary Tree Node

```
57  /**
58   * Returns the left child of this node, or null if it does
59   * not have one.
60   */
61  public LinkedBinaryTreeNode<E> getLeft() {
62      return left;
63  }
64
```


Linked Binary Tree Node

```
65  /**
66   * Inserts child as the
67   * left child of this node. If this node already has a left
68   * child it is removed.
69   * @exception IllegalArgumentException if the child is
70   * an ancestor of this node, since that would make
71   * a cycle in the tree.
72   */
73  public void setLeft(LinkedBinaryTreeNode<E> childNode) {
74
75      // Ensure the child is not an ancestor.
76      for (LinkedBinaryTreeNode<E> n = this; n != null; n = n.parent) {
77          if (n == childNode) {
78              throw new IllegalArgumentException();
79          }
80      }
81
82      // Break old links, then reconnect properly.
83      if (this.left != null) {
84          left.parent = null;
85      }
86      if (childNode != null) {
87          childNode.parent = this;
88      }
89      this.left = childNode;
90  }
91
```

Linked Binary Tree Node

```
92  /**
93   * Returns the right child of this node, or null if it does
94   * not have one.
95   */
96  public LinkedBinaryTreeNode<E> getRight() {
97      return right;
98  }
99
```

Linked Binary Tree Node

```
100- /**
101-  * Inserts it as the
102-  * right child of this node. If this node already has a right
103-  * child it is removed.
104-  * @exception IllegalArgumentException if the child is
105-  * an ancestor of this node, since that would make
106-  * a cycle in the tree.
107-  */
108- public void setRight(LinkedBinaryTreeNode<E> childNode) {
109-     // Ensure the child is not an ancestor.
110-     for (LinkedBinaryTreeNode<E> n = this; n != null; n = n.parent) {
111-         if (n == childNode) {
112-             throw new IllegalArgumentException();
113-         }
114-     }
115-
116-     // Break old links, then reconnect properly.
117-     if (right != null) {
118-         right.parent = null;
119-     }
120-     if (childNode != null) {
121-         childNode.parent = this;
122-     }
123-     this.right = childNode;
124- }
```

Linked Binary Tree Node

```
125
126- /**
127    * Removes this node, and all its descendants, from whatever
128    * tree it is in. Does nothing if this node is a root.
129    */
130- public void removeFromParent() {
131    if (parent != null) {
132        if (parent.left == this) {
133            parent.left = null;
134        } else if (parent.right == this) {
135            parent.right = null;
136        }
137        this.parent = null;
138    }
139 }
140
141- /**
142    * Returns the number of children of this node
143    *
144    */
145- public int numChildren(){
146    int count = 0;
147    if (this.getLeft() != null)
148        count++;
149    if (this.getRight() != null)
150        count++;
151    return count;
152 }
153 }
```

Visitor Interface

```
1
2 /**
3     * Simple visitor interface.
4     * visit can mean anything
5     */
6 public interface Visitor {
7
8     <E> void visit(LinkedBinaryTreeNode<E> Node);
9 }
10
```

Linked Binary Tree

```
1
2 public class LinkedBinaryTree<E> implements Visitor{
3
4     /**
5      * Root of the tree.
6      */
7     protected LinkedBinaryTreeNode<E> root = null;
8
9     /**
10     * Number of nodes in the tree.
11     */
12     private int size = 0;
13
14     /**
15     * constructs an empty binary tree.
16     */
17     public LinkedBinaryTree(){ }
```

Linked Binary Tree

```
18
19  /**
20   * returns the number of nodes in the tree.
21   */
22
23  public int size(){
24      return size;
25  }
26
27  /**
28   * Returns the root position of the tree.
29   */
30
31  public LinkedBinaryTreeNode<E> root(){
32      return root;
33  }
34
35  /**
36   * Returns whether the tree is empty or not.
37   */
38
39  public boolean isEmpty(){
40      return size() == 0;
41  }
```

Linked Binary Tree

```
42
43  /**
44   * Creates a new node storing element e.
45   */
46
47  protected LinkedBinaryTreeNode<E> createNode(E e,
48      LinkedBinaryTreeNode<E> parent, LinkedBinaryTreeNode<E> left,
49      LinkedBinaryTreeNode<E> right ){
50
51      return new LinkedBinaryTreeNode<E>(e, parent, left, right);
52  }
53
```


Linked Binary Tree

```
54
55  /**
56   * Place element e at the root of an empty tree.
57   */
58
59  public LinkedBinaryTreeNode<E> addRoot(E e) throws IllegalStateException {
60      if(!isEmpty()) throw new IllegalStateException("Tree is not empty");
61      root = createNode(e, null, null, null);
62      size = 1;
63      return root;
64  }
```

Linked Binary Tree

```
65
66  /**
67   * Create a new left child of the node n
68   */
69  public LinkedBinaryTreeNode<E> addLeft(LinkedBinaryTreeNode<E> n, E e){
70
71      LinkedBinaryTreeNode<E> child = createNode(e, null, null, null);
72      n.setLeft(child);
73      size++;
74      return child;
75  }
76
```

Linked Binary Tree

```
77  /**
78   * Create a new right child of the node n
79   */
80  public LinkedBinaryTreeNode<E> addRight(LinkedBinaryTreeNode<E> n, E e){
81
82      LinkedBinaryTreeNode<E> child = createNode(e, null, null, null);
83      n.setRight(child);
84      size++;
85      return child;
86  }
87
```

Linked Binary Tree

```
88  /**
89   * Replaces the element at node n with e
90   */
91  public void set(LinkedBinaryTreeNode<E> n, E e){
92      n.setData(e);
93  }
94
```

Linked Binary Tree

```
95  /**
96   * Removes the node n and replaces it with its child if any
97   */
98  public void remove(LinkedBinaryTreeNode<E> n) throws IllegalArgumentException{
99      if (n.numChildren() == 2) throw new IllegalArgumentException("node has two children");
100     LinkedBinaryTreeNode<E> child = (n.getLeft() != null ? n.getLeft() : n.getRight());
101     if (child != null)
102         child.setParent(n.getParent());
103     if (n == root)
104         root = child;
105     else{
106         LinkedBinaryTreeNode<E> parent = n.getParent();
107         if (n == parent.getLeft())
108             parent.setLeft(child);
109         else
110             parent.setRight(child);
111     }
112 }
113
```

Linked Binary Tree

```
114  /**
115   * Visits the nodes in this tree in preorder.
116   */
117  public void traversePreorder() throws IllegalStateException {
118      if (isEmpty()) throw new IllegalStateException("Tree is empty");
119      traversePreorder(root);
120  }
121
122  private void traversePreorder(LinkedBinaryTreeNode<E> node){
123      visit(node);
124      if (node.getLeft() != null) traversePreorder(node.getLeft());
125      if (node.getRight() != null) traversePreorder(node.getRight());
126  }
127
```

Linked Binary Tree

```
128  /**
129   * Visits the nodes in this tree in postorder.
130   */
131  public void traversePostorder() throws IllegalStateException {
132      if (isEmpty()) throw new IllegalStateException("Tree is empty");
133      traversePostorder(root);
134  }
135
136  private void traversePostorder(LinkedBinaryTreeNode<E> node){
137      if (node.getLeft() != null) traversePostorder(node.getLeft());
138      if (node.getRight() != null) traversePostorder(node.getRight());
139      visit(node);
140  }
141
```

Linked Binary Tree

```
142  /**
143   * Visits the nodes in this tree in inorder.
144   */
145  public void traverseInorder() throws IllegalStateException {
146      if (isEmpty()) throw new IllegalStateException("Tree is empty");
147      traverseInorder(root);
148  }
149
150  private void traverseInorder(LinkedBinaryTreeNode<E> node){
151      if (node.getLeft() != null) traverseInorder(node.getLeft());
152      visit(node);
153      if (node.getRight() != null) traverseInorder(node.getRight());
154  }
155
```


Linked Binary Tree

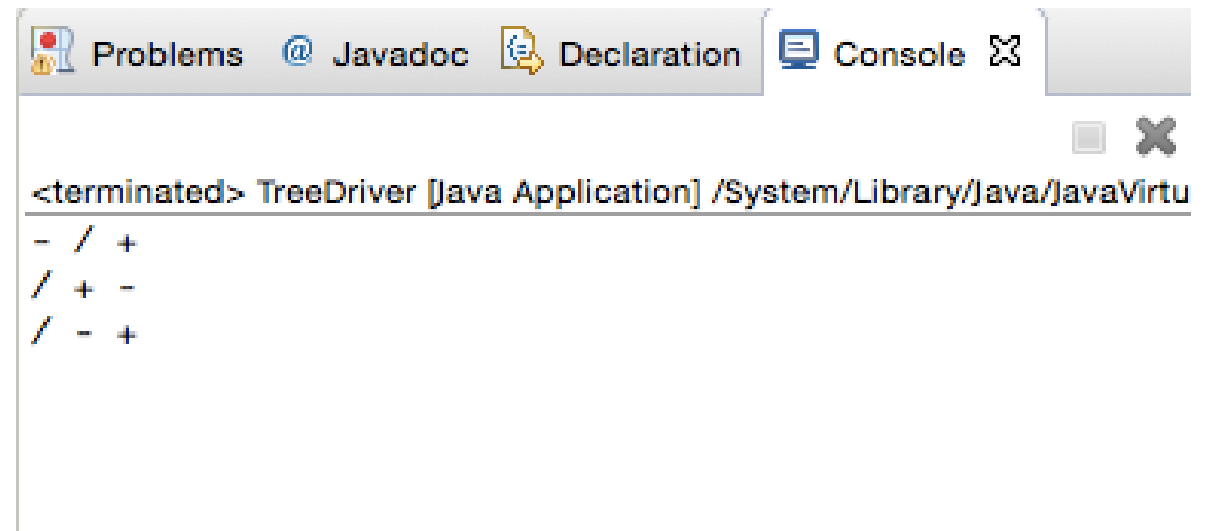
```
156  /**
157     * visiting a node in this case means printing its value.
158     */
159  @Override
160  public <E> void visit(LinkedBinaryTreeNode<E> Node) {
161      // TODO Auto-generated method stub
162      System.out.print(""+Node.getData()+" ");
163
164  }
165 }
166
```

Tree Driver Class

```
1
2 /**
3  * Simple driver class for testing our binary tree implementation.
4  */
5
6 public class TreeDriver {
7
8     /**
9     * Makes an application which adds three string elements to a binary tree.
10    * and performs tree traversals
11    */
12    public static void main(String args[]){
13
14        LinkedBinaryTree<String> tree = new LinkedBinaryTree<String>();
15        LinkedBinaryTreeNode<String> node;
16    }
```

Tree Driver Class

```
17 //adding the root
18 node = tree.addRoot("-");
19 //adding the left child
20 tree.addLeft(node, "/");
21 //adding the right child
22 tree.addRight(node, "+");
23
24 //preorder traversal
25 tree.traversePreorder();
26 System.out.println();
27
28 //postorder traversal
29 tree.traversePostorder();
30 System.out.println();
31
32 //inorder traversal
33 tree.traverseInorder();
34
35
36 }
37 }
```



The screenshot shows an IDE window with tabs for Problems, Javadoc, Declaration, and Console. The Console tab is active, displaying the output of the TreeDriver application. The output consists of three lines: "- / +", "/ + -", and "/ - +".

```
<terminated> TreeDriver [Java Application] /System/Library/Java/JavaVirtu
- / +
/ + -
/ - +
```

- The intuition behind this implementation was to show the working of a linked binary tree in a simplified manner
- For a more thorough implementation, please consult Chapter. 8 of your textbook

Observations

- Did you notice that there was no generic add method in our implementation?
- Why do you think it is so?