Sorting

QUICK SORT
SORT COMPARISON

Quick sort (Hoare sort, 1959)

Divide-and-conquer approach

Comparative sort

Time complexity

- Worst case O(n²)
- Best case O(n log(n))
- Average case T(n log(n))

Space complexity

- Average case O(n log(n))
- Worst case O(n)

Sort principle

1) **Split** array into two parts where right parts dominates on left.

```
\exists (pivot \text{ from i...j}) \forall (m \text{ from i...j})
```

QSort (array, low, high) {

- if (pivot < m) then A[pivot] < A[m]
- if (*pivot* >= m) then A[*pivot*] >= A[m]



QSort(array, low, mid-1); QSort(array, mid+1, high);

```
mid = split(array, low, high);
```

Split algorithm

```
qsort(int a[], int left, int right) {
                   = left:
                                     // set left index
         int l
         int r
                   = right;
                                      // set right index
         int foo = 0;
         int pivot = a[(1 + r) / 2]; // pivot selection is not determined
         while (l <= r) {
                   while ((a[1] < pivot) && (1 <= right)) 1++; // moving while ok
                   while ((a[r] > pivot) \&\& (r >= left)) r--; // moving while ok
                   if (1 <= r) {
                             foo = a[1]; a[1] = a[r]; a[r] = foo; // swap
                             1++;
                             r--;
                                                                  pivot
                                                    9,8,7,6,5,4,3,2,1
         }
         if (r > left) qsort(a, left, r);
         if (1 < right) qsort(a, 1, right);</pre>
                                                    9,8,7,6,5,4,3,2,1
                                                           compare and swap
                                                    1,8,7,6,5,4,3,2,9
```

Drawbacks and best cases

Quick sort:

- Takes O(n) comparisons for <u>array of equal</u> <u>elements</u> (with 3 part splitting)
- Cache-friendly
- Parallelizable
- Works for linked lists (initially streamers)
- Worst case n² for sorted array
 - if each split produces 1 and n-1 element arrays (when selecting 1st element as a pivot in a sorted array)
- Unstable sort

Merge sort:

- Highly parallelizable for large amounts of data
- Worst case is still n log(n)
- Stable sort
- Takes <u>n</u> additional space always

Heap sort:

- O(1) additional memory
- Worst case is still n log(n)
- Cache-unfriendly (jumps over the array)
- Unstable sort