

Computationally Hard Problems

Summary

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Today

- ▶ No new material
- ▶ A very short course summary
- ▶ Some remarks on the exam
- ▶ Example of NP-completeness proof
- ▶ Results of competition
- ▶ Room for questions and discussions

Afterwards: exercise sessions where we continue to work on the old exam. You are welcome to ask further questions then.

What was in the Course

- ▶ Formal stuff: Languages, running times
- ▶ Randomized algorithms: Definition, complexity classes
- ▶ Decision vs. optimization problems
- ▶ \mathcal{NP} -completeness: Definition, proving a problem is in \mathcal{NP} , proving a problem to be \mathcal{NP} -complete, list of \mathcal{NP} -complete problems
- ▶ Using randomization to solve problems exactly, approximately or heuristically

Exam

The exam is on Monday, 14.12.2020, 9:00–13:00 in 101E/Sportshal 1 and 09:00-14:00 in 127/013 (without guarantee; check DTU Inside and watch all announcements).

The exam will consist of seven assignments.

All aids allowed (you may bring your computer with copies of lecture notes etc.)
No internet access.

The lecture are the definite source of reference and comprise the whole curriculum.

Things worth to remember

- ▶ The non-mandatory and mandatory assignments
- ▶ The curriculum of the course (→ multiple-choice exercise)

Structure of exam

- ▶ Design of a language (exam exercise 1)
- ▶ Conversion decision to optimization (exam exercise 2).
- ▶ Proving a problem in \mathcal{NP} (exam exercise 3).
- ▶ Proving a problem \mathcal{NP} -complete (exam exercise 4).
- ▶ Analysis of randomized (approximation) algorithms (exam exercises 5–6).
- ▶ Multiple choice (exam exercise 7)

Some Advice

- ▶ No need to panic.
- ▶ Read the assignment thoroughly.
- ▶ Make sure you know what you are asked to do.
- ▶ Do not answer questions that are not asked.
- ▶ Remember the assignments from the exercise sessions.
- ▶ Write your solutions directly into the answer boxes on the exam set.

Example of Exam Assignment

Exercise 1:



a) Specify the alphabet Σ you use.

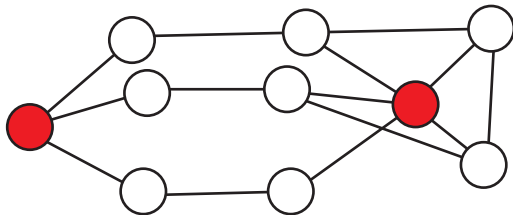
b) Specify how the language L is defined and describe how the above example is coded in your language.

Yet Another \mathcal{NP} -completeness Proof

Problem [NEIGHBORWATCH]

Input: An undirected, connected graph $G = (V, E)$ and a positive integer k .

Output: YES if G has a *guard set* of size at most k , and NO otherwise. A **guard set** is a set $U \subseteq V$ such that for every $v \in V \setminus U$ there is $u \in U$ such that $\{u, v\} \in E$.



Rep.: What to do

In order to prove that a problem P is \mathcal{NP} -complete we have to:

- 1) Prove that $P \in \mathcal{NP}$.
- 2) Find a suitable problem P_c which is known to be \mathcal{NP} -complete.
- 3) Prove $P_c \leq_p P$, especially:
 - 3a) Describe a transformation T which transforms every instance X of P_c into an instance $T(X)$ of P and which runs polynomial in the size $\|X\|$ of X .
 - 3b) Show that if the answer to X is YES then so is the answer to $T(X)$.
 - 3c) Show that if the answer to $T(X)$ is YES then so is the answer to X .

1) NW is in \mathcal{NP}

As always ...

2) NW is \mathcal{NP} -complete: Reference Problem

Find a problem to reduce from.

Any suggestions?

We use VERTEXCOVER.

We show

$$\text{NEIGHBORWATCH} \leq_p \text{VERTEXCOVER}.$$

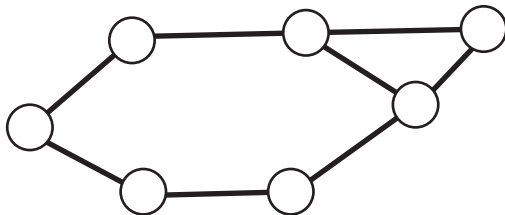
NO!!!!!! We show

$$\text{VERTEXCOVER} \leq_p \text{NEIGHBORWATCH}.$$

3a) The Reduction

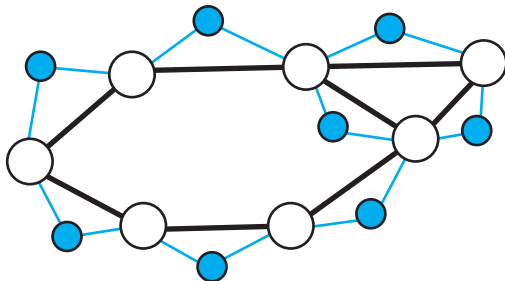
Given an instance $(G = (V, E), k)$ of VC we transform it to an instance $(G' = (V', E'), k)$ of NW as follows. G' is similar to G , with a vertex added for each edge, and connected to both ends of the original edge.

The VC instance



3a) The Reduction

The transformed NW instance



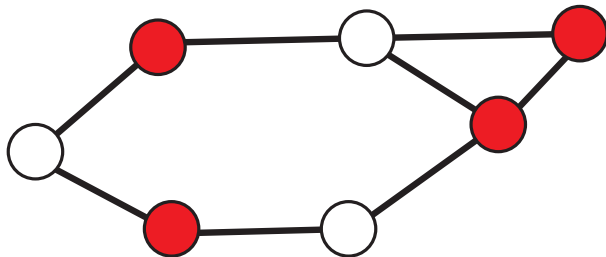
The parameter k is the same as in the VC instance.

This transformation can be done in polynomial time.

3b) \Rightarrow

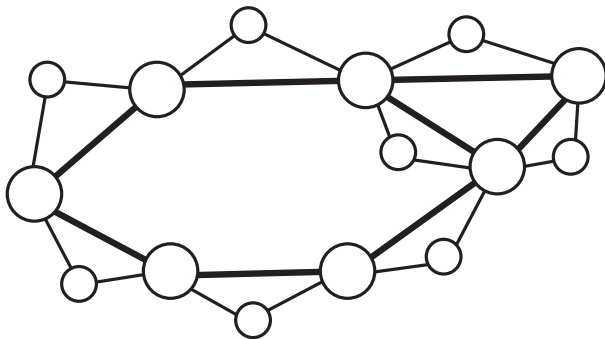
Let U be a vertex cover of G of size at most k . Then we claim that there is also a guard set of size at most k in G' .

The VC instance with vertex cover U (red)



3b) \Rightarrow

U is also a guard set in G' . The NW instance

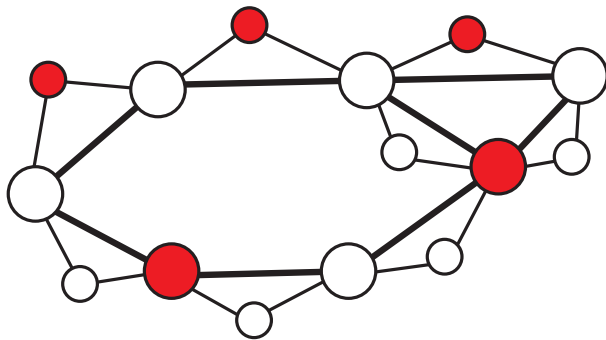


Why is that so?

3c) \Leftarrow

Suppose that U is a guard set in G' of size at most k . Then we claim that there is also a vertex cover of size at most k in G' .

The NW instance

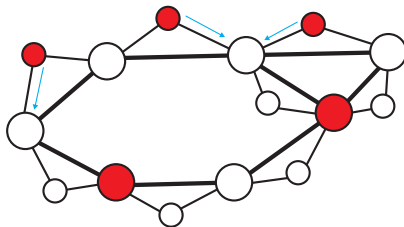


3c) \Leftarrow

Modify guard set for G' into another, no larger guard set for G' which is also a vertex cover for G .

Move the guards from the nodes that were introduced by the transformation to an arbitrary endpoint of the corresponding edge.

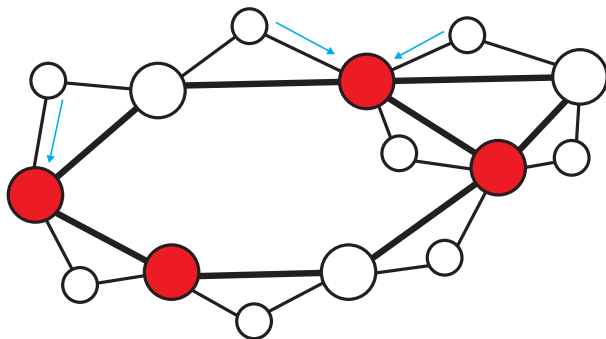
The NW instance



Note: each triangle consisting of original edge and new vertex must contain at least one guard.

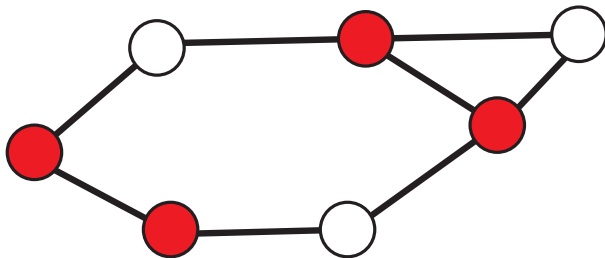
3c) \Leftarrow

The NW instance and the modified guard set



3c) \Leftarrow

The original VC instance and the vertex cover



Competition

Competition

- ▶ 25 teams participated!
- ▶ Score function to be maximized: sum of edge weights minus value of your solution
- ▶ 29 new instances:
 - ▶ Complete graphs with $n = 50, 100, \dots, 250$ and edge weight n
 - ▶ Graphs consisting of path of length n with uniform weights in $[1, 100]$ plus an extra edge of weight $\binom{n}{2}$. Used $n = 1000, 3000, 5000, 8000$.
 - ▶ Graphs consisting of k cliques on c vertices each, joined by $k - 1$ edges; each weight uniform in $[1, 10]$. Used $(100, 2)$, $(100, 4)$, $(100, 6)$, $(100, 8)$, $(100, 10)$ and $(200, 3)$, $(200, 6)$, $(200, 9)$, $(200, 12)$, $(200, 15)$ for (c, k)
 - ▶ Erdős-Rényi random graphs with $n = 100, 150, \dots, 550$, expected $1.1 \binom{n}{2} \ln n$ edges (some of these unconnected) and uniform weights in $[1, n^2]$
- ▶ Time budget was 180 seconds in the first two settings, 60 seconds in the third and 90 seconds in the last setting.
- ▶ Showing the anonymized scoreboard now.

224236476	2199346059	198992585	191700714	191700714			
255426937	248925305	232460180	227262546	227262546	257625670		
124522328	121045739	112432086	108867644	108867644	125456574	125441809	
49000362	48140772	43528742	41204366	41204366	50031636	50013671	5
48884742	47961885	44193441	42789855	42789855	49478233	49401938	4
14352862	13886618	12860638	12472977	12472977	14493266	14501388	1
1299518	1266885	1123394	1135259	1135259	1307330	1303192	
1641011		1640060			1640355		
1311937		1311115			1311425		
981873	974019	981326			981465		
655520	650298	655138	650184	650323			
327372	324864	327383	324906	324947	327436		
270833	266828	269609	266740	266743	270802	270802	
217155	213953	216202	214076	213925	217126	217126	
163355	160947	162756	160960	160855	163348	163348	
108498	106880	107990	106915	106859	108485	108485	
54324	53514	54108	53455	53516	54312		
50	96	50	50	50	99		
39	98	39	39	39	98		
95	97	95	95	95	97		
74	94	74	74	74	94		
7719000	7719000	7719000	7719000	7719000	7719000		
3940200	3940200	3940200	3940200	3940200	3940200		
1653900	1653900	1653900	1653900	1653900	1653900	1653900	
485100	485100	485100	485100	485100	485100	485100	
58800	58800	58800	58800	58800	58800	58800	
737312521 (81.43s)	718872051 (104.97s)	665174111 (83.93s)	641067855 (66.80s)	641067787 (108.73s)	517324871 (109.86s)	242859798 (324.31s)	
		Rikke Bertram Rasmussen (s174493) Anton Stockmarr (s164170)					

224236476	219936059	198992585	191700714	191700714	257625670	125441809
255426937	249925305	232460180	227262546	227262546		50013671
124522328	121045739	112432086	108867644	108867644	125456574	125441809
49000362	48140772	43528742	41204366	41204366	50031636	50013671
48884702	47961885	44193441	42789855	42789855	49478233	49401938
14352862	13986618	12860638	12472977	12472977	14493266	14501388
1299518	1266885	1123394	1135259	1135259	1307330	1303192
1641011		1640060			1640355	
1311937		1311115			1311425	
981873	974019	981326			981465	
655520	650298	655138	650184	650323		
327572	324864	327383	324906	324947	327436	
270833	266828	269609	266740	266743	270802	
217155	213953	216202	214076	213925	217126	
163355	160947	162756	160960	160855	163348	
108498	106880	107990	106915	106859	108485	
54324	53514	54108	53455	53516	54312	
50	96	50	50	50	99	
39	98	39	39	39	98	
95	97	95	95	95	97	
74	94	74	74	74	94	
7719000	7719000	7719000	7719000	7719000	7719000	
3940200	3940200	3940200	3940200	3940200	3940200	
1653900	1653900	1653900	1653900	1653900	1653900	1653900
485100	485100	485100	485100	485100	485100	485100
58800	58800	58800	58800	58800	58800	58800
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Sindri Pétur Ingimundarson (s192586)	Laouen Pablo Killian Fernet (s192612)	Rikke Bertram Rasmussen (s174493)				
Anton Stockmarr (s164170)						

Ask Dörge Kunding (s174266) Asbjørn Pedersen Kaad (s174282)	737312521 (81.43s)	224236476 255426937	219936059 249925305	198992585 232460180	191700714 227262546	191700714 227262546	257625670	
Sindri Pétur Ingimundarson (s192586) Laouen Pablo Killian Fernet (s192612)	718872051 (104.97s)	124522328 49000362 48884742	121045739 48140772 47961885	112432086 43528742 44193441	108867644 41204366 42789855	108867644 41204366 42789855	125456574 50031636 49478233	125441809 50013671 49401938
Rikke Bertram Rasmussen (s174493) Anton Stockmarr (s164170)	665174111 (83.93s)	14352862 1299518 1641011 1311937 981873 655520 327572 270833 217155 163355 108498 54324 50 39 95 74 7719000 3940200 1653900 485100 58800	13986618 1266885 1264060 1311115 981326 650298 324864 266828 213953 160947 106880 53514 96 98 97 94 7719000 3940200 1653900 485100 58800	12860638 1123594 1640060 1311115 981326 655138 327383 269609 216202 162756 107990 54108 50 39 95 74 7719000 3940200 1653900 485100 58800	12472977 1135259 1135259 1135259 1135259 650184 324906 266740 214076 160960 106915 53455 50 39 95 74 7719000 3940200 1653900 485100 58800	12472977 1135259 1135259 1135259 1135259 650323 324947 266743 213925 160955 106859 53516 50 39 95 74 7719000 3940200 1653900 485100 58800	14493266 1307350 1640355 1311425 981465 327436 270802 217126 163348 108485 54312 99 98 97 94 7719000 3940200 1653900 485100 58800	14501388 1303192 1640355 1311425 981465 327436 270802 217126 163348 108485 54312 99 98 97 94 7719000 3940200 1653900 485100 58800
	242859798 (324.31s)							

Questions

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The End

Thank you for the positive course evaluation!

Big thank you to Amir and Andrei.

Good luck with the exam!

And happy holidays!