

Theorem and algorithm of Agrawal, Kayal, Saxena

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"We present a deterministic polynomial time algorithm that determines whether an input number n is prime or composite."

(Abstract of "PRIMES is in P" by Agrawal, Kayal, Saxena)

This talk will cover the algorithm to determine whether an integer is a prime or is composed by Agrawal, Kayal and Saxena and give some background on previous attempts to find such an algorithm.

Tools to check result of primality test

- If n prime all numbers smaller than n are coprime ("teilerfremd") to n
- n prime $\Rightarrow n - 1$ values mod n which are coprime to n
- These values form a cyclic group with multiplication \Rightarrow there is a generator g of this group with $ord(g) = n - 1$
- If g not generator then there is a prime $q \leq n - 1$ with $q | n - 1$ for which $g^{\frac{n-1}{q}} \equiv 1 \pmod{n}$
- with $Q = \{q : q \text{ prime and } q | n - 1\}$ and a generator g one can check a given integer n is prime
- Need to check $g^{n-1} \equiv 1 \pmod{n}$ and $g^{\frac{n-1}{q}} \not\equiv 1 \pmod{n} \forall q \in Q$
- So primality tests are in NP
- This is not algorithm we look for, as we would have to factor $n-1$, only checking is "easy"

Wilson's Theorem (1770)

- integer $n \geq 2$ is prime $\Leftrightarrow n$ divides $(n - 1)! + 1$
- Difficult to convert into algorithm as $(n - 1)!$ is hard to compute

Fermat (1637)

- Known Binomial theorem:
$$(x + y)^n = \sum_{i=0}^n \binom{n}{i} x^i y^{n-i}$$
- In fields K with $char(K) = p$ (p a prime):
$$(x + y)^p = x^p + y^p$$
- Characterisation of primes:
 - p prime $\Rightarrow p | a^p - a \forall a \in \mathbb{Z}$
 - p prime $\Rightarrow a^p \equiv a \pmod{p}$
- Problem: there are also composites for which this holds

Square roots

- 1 has at most 2 square roots in $\mathbb{Z}/p\mathbb{Z}$ (p a prime): 1 and $-1 \equiv p - 1$
- 1 has at least for different square roots in $\mathbb{Z}/n\mathbb{Z}$ (n an integer composed of at least two different primes)
- Call $a \in \mathbb{Z}$ a witness to n if the sequence $a^{n-1} \pmod{n}, a^{\frac{n-1}{2}} \pmod{n}, a^{\frac{n-1}{4}} \pmod{n}, \dots$ Does not reach 1 or -1
- Can show that at least half of the numbers smaller than n are witnesses
- Test numbers to find witness produces „industrial strength prime“

Agrawal, Kayal, Saxena

- Theorem 1:
An integer n is prime if
$$(x + 1)^n \equiv x^n + 1 \pmod{n}$$
 in $\mathbb{Z}[x]$
- **Main theorem:**
- For given integer $n \geq 2$, let r be a positive integer $< n$, for which n has order $> (\log(n))^2 \pmod{r}$. Then n is prime if and only if
 - n is not a perfect power
 - n does not have a prime factor $\leq r$
 - $(x + a)^n \equiv x^n + a \pmod{n, x^r - 1}$ in $\mathbb{Z}[x]$ for each integer a with $1 \leq a \leq \sqrt{r} \log(n)$

Algorithm

- Input: integer n
 - Output: " n is prime" or " n is composite"
1. Determine whether n is perfect power
 2. Find integer r with $ord(n) \pmod{r} > \log(n)^2$
 3. Compute $n^j \pmod{q}$ for $j = 1, \dots, [\log(n)^2]$ and each integer $q > [\log(n)^2]$ until find q for which non of $n^j \equiv 1 \pmod{q}$
 4. Take $r = q$
 5. Determine whether $gcd(a, n) > 1$ for an $a \leq r$
 6. Determine whether $(x + a)^n \equiv x^n + a \pmod{n, x^r - 1}$ for $a = 1, 2, \dots, [\sqrt{r} \log(n)]$