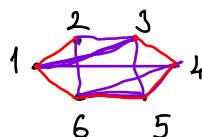


Dogukan Sertbas

## Probabilistic Number Theory

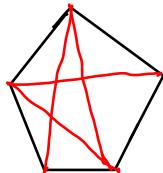
### §0 Prologue (Ramsey Theory)

Example.

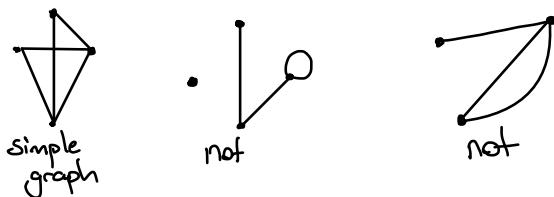


Numbers are person, edges are relationship.

Example.



Definition. A simple graph  $G = (V, E)$  is a tuple, where  $V$  denotes the set of vertices, and  $E$  is the set of edges which consists unordered pairs  $\{x, y\}$  such that  $x, y \in V$  ( $E \subseteq P_2(V) = \{\{x, y\} : x, y \in V\}$ )



Definition. A complete graph  $K_n$  is a simple graph on  $n$  vertices ( $|V| = n$ ) such that  $\forall x, y \in V \quad \{x, y\} \in E$ .

Definition. A two-coloring of the edges  $K_n$  is a map  $\chi: E(K_n) \rightarrow \{\text{Black, White}\}$  which assigns either Black or White to each edge.

? Let  $n > 1$  and  $j, k \in \{1, \dots, n\}$ . For any two-coloring  $\chi$  of  $K_n$ , does  $K_n$  contain either a white monochromatic  $K_j$  or a black monochromatic  $K_k$ ?

$$\hookrightarrow R(j, k) := \min \{n \in \mathbb{N} \mid K_n \text{ contains either a white } K_j \text{ or a black } K_k \text{ for any two coloring of } K_n\} = \min f$$

If  $f$  is nonempty, then Ramsey Theorem  $\forall j, k \in \mathbb{N} \geq 1 \quad R(j, k) < \infty$  and  $R(j, k) \leq \binom{j+k-2}{j-1}$

Proof. We will prove it by induction on  $j+k$ . Note that

$$R(j, 1) = R(1, k) = 1$$

$$1 = R(j, 1) \leq \binom{j-1}{j-1} \text{ or } 1 = R(1, k) = \binom{k-1}{0}$$

Now suppose that the statement is true for  $j+k-1$ . Let  $n = R(j-1, k) + R(j, k-1)$  and  $\chi$  be any two-coloring of  $K_n$ . Take any vertex  $v$  of  $K_n$ .

Define

$$W = \{w, v\} : \chi(\{w, v\}) = \text{white}$$

$$B = \{b, v\} : \chi(\{b, v\}) = \text{black}$$

Observe that  $|W| + |B| = n - 1 = \mathcal{L}(j-1, k) + \mathcal{L}(j, k-1) - 1$ , then  $|W| > \mathcal{L}(j-1, k)$  or  $|B| > \mathcal{L}(j, k-1)$  (+)

(Suppose not,  $|W| \leq \mathcal{L}(j-1, k) \Leftrightarrow |W| \leq \mathcal{L}(j-1, k) - 1$   
 $|B| \leq \mathcal{L}(j, k-1) \Leftrightarrow |B| \leq \mathcal{L}(j, k-1) - 1$ )

It gives  $|W| + |B| \leq \mathcal{L}(j-1, k) + \mathcal{L}(j, k-1) - 2 = n - 2$ .

Case 1. \* holds, there is a white monochromatic  $K_{j-1}$  lies in  $W$  or there is a black monochromatic  $K_k$  lies in  $B$ .

If the second case holds, then a black  $K_k$  lies in  $K_n$ . Otherwise, since all the elements of  $W$  are connected with white edges to  $v$ , this means that there is a monochromatic white  $K_j$  lies  $K_n$  (Similar situation also occurs when (+) holds)

Thus  $\mathcal{L}(j, k) \leq \mathcal{L}(j-1, k) + \mathcal{L}(j, k-1)$

$$\begin{aligned} &\leq \binom{j-1+k-2}{k-1} + \binom{j+k-3}{k-2} \\ &= \binom{j+k-2}{k-1} \end{aligned}$$

□

$$\mathcal{L}(2, 2) = 2$$

$$\mathcal{L}(4, 4) = 18$$

$$102 \leq \mathcal{L}(6, 6) \leq 165$$

$$\mathcal{L}(3, 3) = 6$$

$$43 \leq \mathcal{L}(5, 5) \leq 49$$

Remark. For any  $k \geq 1$ , we have  $\mathcal{L}(k, k) \leq \binom{k+k-2}{k-1} = \binom{2k-2}{k-1} \leq \sum_{i=0}^{2k-2} \binom{2k-2}{i} i!^{-2k-2-i}$

? What can we say about lower case bound for  $\mathcal{L}(k, k)$  depending on  $k$ ?

## § 1 Basic Probability Theory.

### 1.1 Probability Spaces.

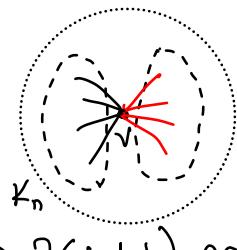
Definition. Let  $X$  be non-empty set and  $\mathcal{A}$  be a collection of subsets of  $X$  ( $\mathcal{A} \subseteq \mathcal{P}(X)$ ).  $\mathcal{A}$  is called  $\sigma$ -algebra if

i.  $\emptyset \in \mathcal{A}$

ex.  $\mathcal{A} = \mathcal{P}(X)$ ,  $A = \{\emptyset, X\}$

ii.  $\forall A \in \mathcal{A}, A^c \in \mathcal{A}$

iii.  $\{\mathbb{A}_i\}_{i=1}^{\infty} \subseteq \mathcal{A} \Rightarrow \bigcup_{i=1}^{\infty} A_i \in \mathcal{A}$



**Definition.** A probability space  $(\Omega, \mathcal{F}, P_r)$  is a tuple where  $\Omega$  is nonempty set which stands for the set of possible outcomes  $\mathcal{F}$  is a  $\sigma$ -algebra on  $\Omega$  where  $\mathcal{F}$  contains the all possible events and  $P_r: \mathcal{F} \rightarrow [0,1]$  is a function such that

- $P_r(\Omega) = 1$

- For any sequence  $\{\mathfrak{A}_i\}_{i=1}^{\infty} \subseteq \mathcal{F}$  of disjoint sets, we have  $P_r\left(\bigcup_{i=1}^{\infty} \mathfrak{A}_i\right) = \sum_{i=1}^{\infty} P_r(\mathfrak{A}_i)$  (countable additivity)

**Example. (tossing a coin)**  $\Omega = \{\text{H, T}\}$ ,  $\mathcal{F} = \mathcal{P}(\Omega) = \{\emptyset, \{\text{H}\}, \{\text{T}\}, \{\text{H, T}\}\}$ . If we flip a fair coin, then  $P_r(\{\text{H}\}) = P_r(\{\text{T}\}) = 1/2$ .  $P_r: \mathcal{F} \rightarrow [0,1]$  where  $\emptyset \mapsto 0, \{\text{H}\} \mapsto 1/2, \{\text{T}\} \mapsto 1/2, \Omega \mapsto 1$ .

**Example. (throw 2 dice)**  $\Omega = \{(i,j) : 1 \leq i, j \leq 6\}$ . If  $\mathfrak{A} \in \mathcal{F}$  represents the event that of the faces of two dice is  $\mathfrak{A}$ , then

$$\mathfrak{A} = \{(1,6), (2,5), (3,4), \dots, (6,1)\}$$

If we deal fair dice, then  $P_r(\mathfrak{A}) = \frac{|\mathfrak{A}|}{36} = \frac{6}{36} = 1/6$

## Properties

P1.  $\forall \mathfrak{E} \in \mathcal{F}, P_r(\mathfrak{E}^c) = 1 - P_r(\mathfrak{E})$ .

P2.  $P_r(\emptyset) = 0$ .

P3.  $\forall \mathfrak{E}, \mathfrak{F} \in \mathcal{F}, \mathfrak{E} \subseteq \mathfrak{F}$ , then  $P_r(\mathfrak{E}) \leq P_r(\mathfrak{F})$  ( $\mathcal{F} = (\mathfrak{F} \setminus \mathfrak{E}) \cup \mathfrak{E}$ )

P4. (Inclusion-Exclusion Property).

$$P_r(\mathfrak{E} \cup \mathfrak{F}) = P_r(\mathfrak{E}) + P_r(\mathfrak{F}) - P_r(\mathfrak{E} \cap \mathfrak{F}) \quad (\mathfrak{E} \cup \mathfrak{F} = \mathfrak{E} \setminus \mathfrak{F} \cup (\mathfrak{F} \setminus \mathfrak{E}) \cup (\mathfrak{E} \cap \mathfrak{F}))$$

P5. (Boole's Inequality) If  $\{\mathfrak{A}_i\}_{i=1}^{\infty} \subseteq \mathcal{F}$ , then  $P_r\left(\bigcup_{i=1}^{\infty} \mathfrak{A}_i\right) \leq \sum_{i=1}^{\infty} P_r(\mathfrak{A}_i)$ .

**Proof.** Define  $B_1 := A_1$  and  $B_k := A_k \setminus \bigcup_{i=1}^{k-1} A_i$ . Observe that  $\forall k \in \mathbb{N}$ ,  $B_k \subseteq A_k$ . Also

$\bigcup_{i=1}^{\infty} B_i = \bigcup_{i=1}^{\infty} A_i$ . Note that all  $B_i$ 's are disjoint. Hence,

$$P_r\left(\bigcup_{i=1}^{\infty} A_i\right) = P_r\left(\bigcup_{i=1}^{\infty} B_i\right) = \sum_{i=1}^{\infty} P_r(B_i) \leq \sum_{i=1}^{\infty} P_r(A_i)$$

□

**Corollary 1.** Let  $\{\mathfrak{A}_i\}_{i=1}^{\infty} \subseteq \mathcal{F}$  be a sequence of events. If  $\sum_{i=1}^{\infty} P_r(A_i) < 1$ , then  $P_r\left(\bigcap_{i=1}^{\infty} A_i^c\right) > 0$ . In other words, there is an event where none of  $A_i$ 's occur.

**Proof (sketch).**

$$P_r\left(\bigcap_{i=1}^{\infty} A_i^c\right) = P_r\left(\left(\bigcup_{i=1}^{\infty} A_i\right)^c\right) = 1 - P_r\left(\bigcup_{i=1}^{\infty} A_i\right) \geq 1 - \sum_{i=1}^{\infty} P_r(A_i) > 0.$$

## Probabilistic Method.

**Main Idea.** To show the occurrence of an event  $\mathcal{E}$ , prove that  $\Pr(\mathcal{E}) > 0$

$\hookrightarrow$  (Diagonal) Ramsey Numbers.

**Theorem 1.** For any  $k \in \mathbb{N} \geq 2$ , let  $r(k) = \max \{n \in \mathbb{N} \mid \binom{n}{k} 2^{1-\frac{k}{2}} < 1\}$ . Then  $R(k, k) > r(k)$ . In particular,  $R(k, k) > \frac{k \cdot 2^{\frac{k}{2}}}{2e}$

**Proof.** (Erdős, '47) Remember that Boole's Sieve:

$$\sum_{i \geq 1} \Pr(A_i) < 1 \text{ implies that } \Pr\left(\bigcap_{i \geq 1} A_i^c\right) > 0.$$

for a given  $n \leq r(k)$ , let  $(\Omega, \mathcal{F}, \Pr)$  be a probability space  $\Omega$  contains all possible two colorings of  $K_n$ . Assume that for any coloring  $X \in \Omega$ , we have

$$\Pr(\{X\}) = 1/\binom{n}{2}.$$

Remember that  $\# \mathcal{D}(K_n) = \binom{n}{2}$ . Let  $S$  be any fixed set of vertices in  $K_n$ , where  $|S| = k$ . If  $A_S$  denotes the event that  $S$  constitutes a monochromatic  $K_k$ , then

$$\Pr(A_S) = \underbrace{\prod_{i \in E(S)} \frac{1}{2}}_{\text{being white}} + \underbrace{\prod_{i \in E(S)} \frac{1}{2}}_{\text{being black}} = 2^{1-\frac{k}{2}}$$

Note that

$$\Pr\left(\bigcup_{\substack{S \subseteq V(K_n) \\ |S|=k}} A_S\right) \leq \sum_{\substack{S \subseteq V(K_n) \\ |S|=k}} \Pr(A_S) = \sum_{\substack{S \subseteq V(K_n) \\ |S|=k}} 2^{1-\frac{k}{2}} = 2^{1-\frac{k}{2}} \sum_{\substack{S \subseteq V(K_n) \\ |S|=k}} 1 = \binom{n}{k} 2^{1-\frac{k}{2}}$$

Since  $n \leq r(k)$ , we have  $\binom{n}{k} 2^{1-\frac{k}{2}} < 1$ . By previous Corollary, there exists a two coloring that none of  $K$ 's are monochromatic. Hence  $R(k, k) > r(k)$  by definition.

To find a precise lower bound observe that

$$\binom{n}{k} = \frac{n \cdot (n-1) \cdots (n-k+1)}{k!} < \frac{n^k}{k!}$$

$$\text{Also, } \log(k!) = \log\left(\prod_{j=1}^k j\right) = \sum_{j=1}^k \log j > \int_1^k \log t dt = t \log t - t \Big|_1^k = k \log k - (k-1)$$

$$> k \log k - k$$

This implies  $k! = e^{\log k!} = \left(\frac{k}{e}\right)^k$ . Hence  $\binom{n}{k} 2^{1-\frac{k}{2}} < \frac{n^k e^k}{k^k} \cdot \frac{2}{2^{\frac{k(k-1)}{2}}} = \frac{n^k e^k}{k^k} \cdot \frac{2}{2^{\frac{k(k-1)}{2}}}$ . So if  $\frac{n^k e^k}{k^k} \cdot \frac{2}{2^{\frac{k(k-1)}{2}}} \leq 1$ , or equivalently  $n^k \leq \frac{2^{\frac{k(k-1)}{2}} \cdot k^k}{e^k}$ . We see that there is a two-coloring of  $K$ , which does not contain any monochromatic  $K_k$ . As  $k \geq 2$

$$\frac{k \cdot 2^{\frac{k}{2}}}{2e} \leq \frac{k \cdot 2^{\frac{k}{2}}}{2^{\frac{k(k-1)}{2}} \cdot e^{\frac{k-1}{2}}}.$$

$$\text{Thus, if } n \leq \frac{k \cdot e^{\frac{k-1}{2}}}{2e}, \text{ then } n^k \leq \frac{k^k \cdot 2^{\frac{k(k-1)}{2}}}{2^{\frac{k(k-1)}{2}} \cdot e^k} \left(\binom{n}{k} 2^{1-\frac{k}{2}}\right) < 1$$

Then,

$$R(k, k) > \frac{k \cdot 2^{\frac{k}{2}}}{2e}$$

□

Best Known Bounds for  $R(k, k)$  (Stirling's formula :  $n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$ )

$$(1 + o(1)) \frac{k\sqrt{2}}{e} 2^{\frac{k}{2}} \leq R(k, k)$$

where  $o(1)$  denotes a function  $f(k)$  such that  $\lim_{k \rightarrow \infty} f(k) = 0$ . Recently,  $\exists \varepsilon > 0$  such that

$$R(k, k) < (4 - \varepsilon)^k \quad \left(\frac{2022}{2023} \text{ by I.I.P.}\right)$$

**Remark.** Note that there are  $2^{\binom{n}{2}}$  many possible two-colorings of  $K_n$ . To obtain a monochromatic  $K_k$  in  $K_n$ , we need to count the number of possible colorings which contains a monochromatic  $K_k$ ,

$$\binom{n}{k} \cdot 2 \cdot 2^{\binom{k}{2} - \binom{k}{2}}$$

↓ choosing  $k$ -vertices    ↗ choosing color    → color the rest

If  $\binom{n}{k} 2^{\binom{k}{2} - \binom{k}{2}} < 2^{\binom{n}{2}}$ , then there exists a two-coloring of  $K_n$  where it does not contain any monochromatic  $K_k$ . Hence, if

$$\binom{n}{k} 2^{1 - \binom{k}{2}} < 1$$

then there is a coloring where  $K_n$  does not contain any monochromatic  $K_k$ .

## 1.2 Conditional Probability and Independence

**Definition.** Let  $(\Omega, \mathcal{F}, \Pr)$  be a probability space, and  $A, B \in \mathcal{F}$  two events such that  $\Pr(B) > 0$ . We define

$$\Pr(A|B) = \frac{\Pr(A \cap B)}{\Pr(B)}$$

as the probability of  $A$  given that  $B$  has occurred

### [H-W]. Remark.

1. Notice that for any  $H \in \mathcal{F}$  w/  $\Pr(H) > 0$ , the function

$$\Pr_H : \mathcal{F} \rightarrow [0, 1] \text{ defined by } A \mapsto \Pr_H(A|H)$$

is also a probability measure on  $(\Omega, \mathcal{F}, \Pr_H)$

2.  $\forall H \in \mathcal{F}$ ,  $\Pr(A^c|H) = 1 - \Pr(A|H)$

## The Chain Rule for the Conditional Probabilities

For any  $n \geq 2$ , let  $A_1, \dots, A_n$  be the events whose intersection has strictly positive probability. Then

$$\begin{aligned} \Pr(A_1 \cap \dots \cap A_n) &= \Pr(A_n | A_{n-1} \cap \dots \cap A_1) \\ &\quad \cdot \Pr(A_{n-1} | A_{n-2} \cap \dots \cap A_1) \\ &\quad \cdot \Pr(A_2 | A_1) \cdot \Pr(A_1) \\ &= \prod_{i=1}^n \Pr(A_i | \bigcap_{j=1}^{i-1} A_j) \end{aligned}$$

Proof. Note that  $A_1 \cap \dots \cap A_n \subseteq A_1 \cap \dots \cap A_{n-1} \subseteq \dots \subseteq A_1$ . Hence

$$0 < \Pr\left(\bigcap_{i=1}^n A_i\right) \leq \Pr\left(\bigcap_{i=1}^{n-1} A_i\right) \leq \dots \leq \Pr(A_1)$$

By induction, the rule follows. [H.W]

## The Total Probability Law.

Let  $(\Omega, \mathcal{F}, \Pr)$  be a probability space and  $\{A_i\}_{i=1}^k \subseteq \mathcal{F}$  a partition of  $\Omega$  where  $\Pr(A_i) > 0$  for each  $i \in \mathbb{N}$  (and possibly  $k = \infty$ ) (i.e.  $\Omega = \bigcup A_i$ ). Then  $\forall H \in \mathcal{F}$ ,  $\Pr(H) = \sum_{i=1}^k \Pr(A_i \cap H) = \sum_{i=1}^k \Pr(H | A_i) \cdot \Pr(A_i)$

**Definition.** Let  $A, B \in \mathcal{F}$  be two events in the probability space  $(\Omega, \mathcal{F}, \Pr)$ . These events are called independent if

$$\Pr(A \cap B) = \Pr(A) \Pr(B)$$

### Remark.

- 1) If  $A$  and  $B$  are independent w/  $\Pr(B) > 0$ , then  $\Pr(A|B) = \Pr(A)$
- 2) If  $A$  and  $B$  are independent,  $A$  and  $B^c$  are independent

**Definition.** Let  $\{A_i\}_{i=1}^n \subseteq \mathcal{F}$  be events. We say that  $A_1, \dots, A_n$  are pairwise independent if  $\forall i, j \in \{1, \dots, n\}$   $i \neq j$  implies that  $A_i$  and  $A_j$  are independent. Also  $A_1, \dots, A_n$  are mutually independent. If for any subset  $K \subseteq \{1, \dots, n\}$

$$\Pr\left(\bigcap_{k \in K} A_k\right) = \prod_{k \in K} \Pr(A_k)$$

Moreover, we say that  $A$  is independent of the events  $\{A_j | j \in I\}$  for some  $I \subseteq \{1, \dots, n\}$  if  $\forall K \subseteq I$

$$\Pr(A \cap \bigcap_{k \in K} A_k) = \Pr(A) \cdot \Pr\left(\bigcap_{k \in K} A_k\right)$$

**Remark.** If  $A, B, C$  are independent events, then  $C^c$  are pairwise independent with  $A, B$  and  $A \cap B$ , so  $A^c, B^c, C^c$  are independent. In general,  $A_1, \dots, A_n$  are independent implies that  $A_1^c, \dots, A_n^c$  are independent.

**Proposition** If  $A_1, \dots, A_n$  are independent and  $\Pr(A_i) < 1$  for each  $i \in \{1, \dots, n\}$  then  $\Pr(\bigcap_{i=1}^n A_i^c) > 0$ .

**Proof.** If  $A_1, \dots, A_n$  are independent, then so  $A_1^c, \dots, A_n^c$ . Thus

$$\Pr(\bigcap_{i=1}^n A_i^c) = \prod_{i=1}^n \Pr(A_i^c) = \prod_{i=1}^n (1 - \Pr(A_i)) > 0 \text{ as } \Pr(A_i) \in [0, 1].$$

**Lovasz's Local Lemma.** (The Theory of Probability - S.S. Verbitskiy)

$A_1, \dots, A_n$  events where each  $A_i$ 's mutually independent except at most  $d$  of  $\{A_1, \dots, A_n\} \setminus \{A_i\}$

$$\forall i \in \{1, \dots, n\} \quad \Pr(A_i) \leq p. \quad 4pd \leq 1 \Rightarrow \Pr(\bigcap_{i=1}^n A_i^c) > 0.$$

### § Van der Waerden's Theorem

For any given positive integers  $r, k \geq 1$  there is a natural  $N = N(r, k)$  s.t. if we color the set of integers  $\{1, \dots, N\}$  with  $r$  colors, there is an arithmetic progression of length  $k$  whose elements are of the same color.

$$(X : \{1, \dots, N\} \longrightarrow \{1, \dots, r\} \text{ or } X(a))$$

We define the van der Waerden number  $vdl(r, k)$  as the smallest  $N$  that satisfies vdl theorem

**Theorem (First Lower Bound)** For any  $n \geq 2$  and  $k \geq 2$  we have  $vdl(r, k) \geq \sqrt{2} r^{\frac{k-1}{2}}$

**Proof.** Let  $n \geq \sqrt{2} \cdot r^{\frac{k-1}{2}}$  and  $(\Omega, \mathcal{F}, \Pr)$  be a probability space where  $\Omega$  consists of all possible colorings of  $\{1, \dots, n\}$ . Assume that any coloring occurs equally likely, i.e.,

$$\forall x \in \Omega \quad \Pr(\{x\}) = \left(\frac{1}{r}\right)^n$$

as coloring each number is independent of coloring the others.

Let  $S$  be any  $k$ -term AP in  $\{1, \dots, n\}$ . Denote  $A_S$  the event that  $S$  is monochromatic. Then

$$\Pr(A_S) = r^{1-k}$$

Also observe that the structure of an AP is determined by its first two terms

(Since if we know  $\alpha$  and  $q$ , then we obtain  $a, q$ ) Therefore, there are at most  $\binom{n}{k}$  many possible choices for  $\alpha$  and  $a+q$ . This means that the number of  $k$ -AP in  $\{1, \dots, n\}$  is bounded by  $\binom{n}{k}$ .

$$\sum_{\substack{S \text{ is} \\ a k\text{-AP}}} \Pr_r(A_S) = \sum_S r^{1-k} = \binom{n}{k} r^{1-k} < 1.$$

This implies that there is a coloring that  $\{1, \dots, n\}$  does not contain any number  $k$ -AP

$$\frac{n \cdot (n-1)}{2} r^{k-1} \leq \frac{n^2}{2} r^{k-1} < 1 \Rightarrow n^2 \leq 2 r^{k-1}$$

Since  $n \leq \sqrt{2} r^{\frac{k-1}{2}}$ ,  $\text{VC}(r, k) \neq n$ . Hence  $\text{VC}(r, k) \geq \sqrt{2} r^{\frac{k-1}{2}}$ .

□

$r^k$	1	2	3	$\dots$	7
2	1	3	9	Known	///
3	1	4	27	/ / /	/
4	1	5	76	/ / /	/
.	1	/ /	/ / /	/	/
.					

→ Not Known

Best Known Bounds  $r, k \geq 2 \quad \frac{r^k}{e^{rk}} (1 + o_r(1)) \leq \text{VC}(r, k)$ .

using an improved version of Lovasz Local Lemma

$$\text{VC}(r, k) \leq 2^2 r^{2^{k+q}} \quad (\text{Gower '08}).$$

Theorem (Improved Lower Bounds for  $\text{VC}(r, k)$ ) For all  $r, k \geq 2 \quad \text{VC}(r, k) > \frac{r^k}{4^{rk}}$

Proof. We use Lovasz' Local Lemma and from the previous proof,

$$\Pr_r(A_S) = r^{1-k}$$

for each  $k$ -AP. So, take  $p = r^{1-k}$ .

Also let  $S$  and  $T$  be two different  $k$ -AP in  $\{1, \dots, n\}$ . Notice that  $A_S$  and  $A_T$  are independent, if they do not share any common integer.

Now, we are going to count how many  $k$ -AP's can intersect with  $S$  fix any integer  $x \in S$  and a natural number  $q \geq 0$ . Assume that  $T$  intersects with  $S$  at  $x$  and the difference between the consecutive terms of  $T$  is  $q$ . Observe that there are at most  $k$  such  $T$ 's that intersects with  $x$  (with <sup>the</sup> common difference  $q$ )

Also there are at most  $\lfloor \frac{n}{k} \rfloor$  many choice for  $q$ .

$$1 \leq a \leq a+q \leq \dots \leq a+(k-1)q \leq n \quad \xrightarrow{\text{HW}} \quad q \leq \lfloor \frac{n}{k} \rfloor$$

Finally, there are  $k$  many different  $x$ 's. So we get that there are at most  $nk$  many different

$T$ 's in  $\{1, \dots, n\}$  for which  $A_S$  and  $A_T$  may be dependent, since

$$k \cdot \left\lfloor \frac{n}{k} \right\rfloor \cdot k \stackrel{\# \text{ of different } x_i}{\leq} \frac{n}{k} \cdot k^2 = nk$$

$\xrightarrow{\# \text{ of different } T \text{'s where } x_i \text{ is fixed}}$   $\xrightarrow{\# \text{ of different } q \text{'s}}$

If we take  $d = nk$ , by Lovasz's Local Lemma we obtain that

$$4pd = 4r^{1-k} nk \leq 1$$

implies there is a coloring where none of  $k$ -AP's are monochromatic.

$$(x) \Leftrightarrow n \leq \frac{r^{k-1}}{4k}$$

$$\text{Thus } \chi_1(r, k) > \frac{r^k}{4k}$$

□

### Some Known Bounds

1.  $\chi_1(r, k) \geq 2^k/k^c$  for sufficiently large  $k$  (Szabo '90)

2.  $\chi_1(2, p+1) \geq p \cdot 2^p$  for  $p \in \mathbb{P}$  (Berlekamp '68)

## § Expectation and First Moment Method

**Definition.** A random variable  $(\Omega, \mathcal{F}, P)$   $X$  is a map  $X: \Omega \rightarrow \mathbb{R}$  on a probability space  $(\Omega, \mathcal{F}, P)$  where for each Borel set of reals  $E$  we have  $X^{-1}(E) \in \mathcal{F}$ . (Here Borel set corresponds to an element of the Borel  $\sigma$ -algebra of reals which is the smallest  $\sigma$ -algebra containing the standard topology of reals.)

Here we denote

$$Pr(X \in S) := Pr(\{\omega \in \Omega : X(\omega) \in S\})$$

### Example. (Flip Coin)

$$Pr(X=0) = Pr(X=1) = \frac{1}{2}$$

### Example (uniform r.v.)

$$\forall i \in \{1, \dots, n\} \quad Pr(U=i) = 1/n$$

### Example. (Indicator function as an r.v.)

Let  $(\Omega, \mathcal{F}, P)$  be a probability space and  $A \in \mathcal{F}$ . Then the indicator function

$1_A: \Omega \rightarrow \mathbb{R}$  defined as  $\omega \mapsto \begin{cases} 1, & \text{if } \omega \in A \\ 0, & \text{else} \end{cases}$  is an r.v. (H.W. show it)

If  $|\Omega| = n < \infty$ , the elements of  $A \subseteq \Omega$  are chosen uniformly at random

then we have

$$Pr(A) = Pr\left(\bigcup_{a \in A} \{\omega\}\right) = \sum_{a \in A} Pr(\{\omega\}) = \sum_{a \in A} \frac{1}{n} = \frac{1}{|\Omega|} \sum_{a \in A} 1 = \frac{|A|}{|\Omega|}$$

## Independent Random Variables Set

Let  $(\Omega, \mathcal{F}, P)$  be a probability space  $X: \Omega \rightarrow \mathbb{R}$  and  $Y: \Omega \rightarrow \mathbb{R}$  r.v.s. For some given Borel set  $A, B \subseteq \mathbb{R}$ , we may want to consider the prob. of  $X \in A$  given that  $Y \in B$  where  $P(Y \in B) > 0$ .

In that case

$$P(X \in A | Y \in B) = \frac{P(X \in A, Y \in B)}{P(Y \in B)}$$

where  $P(X \in A, Y \in B) = P(\{w \in \Omega : X(w) \in A, Y(w) \in B\})$

**Definition.** Let  $(\Omega, \mathcal{F}, P)$  be a probability space.  $X, Y: \Omega \rightarrow \mathbb{R}$  r.v.s.

We say that  $X$  and  $Y$  are (pairwise) independent if for any (loc) sets  $A, B \subseteq \mathbb{R}$

$$P(X \in A, Y \in B) = P(X \in A) P(Y \in B).$$

**Definition.** Let  $(\Omega, \mathcal{F}, P)$  be a probability space and  $X$  be a discrete r.v. where  $|Dom(X)| < \infty$ . The expected value or expectation of  $X$  is defined by

$$\mathbb{E}(X) = \sum_{x \in Dom(X)} x \cdot P(X=x).$$

Measure Theory'de  $\int_{\Omega} d\mu$ .

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~~✓~~ Missed one class.

## §4. Variance & Second Moment Method.

### Variance and Covariance.

**Definition.** Let  $(\Omega, \mathcal{F}, P)$  be a probability space and  $X$  a r.v. on  $\Omega$ . Then the variance of  $X$  is defined as

$$\text{Var}(X) := \mathbb{E}((X - \mathbb{E}(X))^2)$$

**Remark.**  $\text{Var}(X) \geq 0$  and standard deviation  $\text{SD}(X) = \sqrt{\text{Var}(X)}$ .

### Properties.

- $\text{Var}(X) = 0 \iff \Pr(X = \mathbb{E}(X)) = 1$
- $\text{Var}(X) = \sum_{\substack{x \in \text{Ran}(X) \\ x = \mathbb{E}(X)}} ((x - \mathbb{E}(X))^2) \Pr(X = x) + \sum_{\substack{x \in \text{Ran}(X) \\ x \neq \mathbb{E}(X)}} (x - \mathbb{E}(X))^2 \Pr(X = x) = 0$ .

**H.W:** Consider the situation  $\Pr(X \neq \mathbb{E}(X)) > 0$ .

- $\forall a, b \in \mathbb{R}$   $\text{Var}(aX + b) = a^2 \text{Var}(X)$ . Consider the

$$\begin{aligned} \text{Var}(aX + b) &= \mathbb{E}((aX + b) - \mathbb{E}(aX + b))^2 \\ &= a^2 (\mathbb{E}(X - \mathbb{E}(X))). \end{aligned}$$

- Chebyshev's Inequality: Let  $X$  be an r.v. and  $\lambda \in \mathbb{R}_{>0}$ . Then

$$\Pr(|X - \mathbb{E}(X)| \geq \lambda \sqrt{\text{Var}(X)}) \leq 1/\lambda^2. \quad (*)$$

**Proof.** First, consider the case when  $\text{Var}(X) = 0$ . This holds iff  $X = \mathbb{E}(X)$  and prob 1 (by  $\text{V}(1)$ ). So,

$$\Pr(|X - \mathbb{E}(X)| > 0) = 1 - \Pr(|X - \mathbb{E}(X)| \leq 0) = 1 - 1 = 0 < 1/\lambda.$$

Now assume that  $\text{Var}(X) > 0$ . Define a new variable  $Y = (X - \mathbb{E}(X))^2$ . By Markov's inequality,  $\Pr(Y > x) \leq \Pr(Y > \alpha) \leq \frac{\mathbb{E}(Y)}{\alpha}$

$$\begin{aligned} \text{Choose } \alpha = \lambda^2 \text{Var}(X) > 0. \text{ Then } \Pr((X - \mathbb{E}(X))^2 > \lambda^2 \text{Var}(X)) &\leq \frac{\mathbb{E}((X - \mathbb{E}(X))^2)}{\lambda^2 \text{Var}(X)} \\ \Leftrightarrow \Pr(|X - \mathbb{E}(X)| > \lambda \sqrt{\text{Var}(X)}) &\leq 1/\lambda^2. \end{aligned}$$

**Remark.**  $(*)$  asserts that  $X = \mathbb{E}(X) + O((\text{Var}(X))^{1/2})$  occurs with high probability.

**Definition.** Let  $X$  and  $Y$  be two r.v. on the same probability space  $(\Omega, \mathcal{F}, P)$ . The covariance of  $X$  and  $Y$  is defined as  $\text{Cov}(X, Y) = \mathbb{E}((X - \mathbb{E}(X))(Y - \mathbb{E}(Y)))$

$$= \mathbb{E}(XY) - \mathbb{E}(X)\mathbb{E}(Y).$$

**Proposition.** (C.1)  $\text{Cov}(X, X) = \text{Var}(X)$

$$(C.2) \text{Cov}(X, Y) = \text{Cov}(Y, X).$$

(C.3) If  $X$  and  $Y$  are independent r.v.s., then  $\text{Cov}(X, Y) = 0$ .

$$(C.4) \text{Var}\left(\sum_{i=1}^n X_i\right) = \sum_{i=1}^n \text{Var}(X_i) + \sum_{\substack{1 \leq i, j \leq n \\ i \neq j}} \text{Cov}(X_i, X_j).$$

↳ (First compute  $\text{Var}(X+Y)$ , then do ...)

(V.5) If  $\{X_i\}_{i=1}^n$  is a set of pairwise independent r.v.s., then

$$\text{Var}\left(\sum_{i=1}^n X_i\right) = \sum_{i=1}^n \text{Var}(X_i)$$

(V.6) Assume that  $B$  is some randomly generated subset of  $A$ . Then

$$\text{Var}(1_B) = \sum_{a \in A} \Pr(a \in B) \Pr(a \in B)^2 + \sum_{\substack{a, a' \in A \\ a \neq a'}} \text{Cov}(1_{a \in B}, 1_{a' \in B})$$

Book: Tao & Vu - Additive Combinatorics.

In particular, if the events  $a \in B$  are pairwise independent, then

$$\text{Var}(1_B) = \sum_{a \in A} (\Pr(a \in B) - \Pr(a \in B)^2)$$

Book: Coppersmith & Morty - An Introduction to Sieve Theory → lecture Notes.

Hardy-Ramanujan Th.

$$\begin{aligned} w(n) &= \sum_{p|n} 1 - \# \text{ of prime divisors of } n. \\ &= \sum_{p \leq n} \frac{1}{p} \end{aligned}$$

Theorem. Let  $f(n)$  tend to infinity arbitrarily slowly. Then

$$\#\{k \leq n : |w(k) - \log \log k| > f(n) \sqrt{\log \log k}\} \sim O(n).$$

Erdős-Kac Theorem ('40) Let  $\lambda \in \mathbb{R}$  be fixed. Then

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{1}{n} |\{k \in [1, n] : w(k) > \log \log n + \lambda \sqrt{\log \log n}\}| \\ = \int_{-\infty}^{\infty} \frac{e^{-t^2/2}}{\sqrt{2\pi}} dt \end{aligned}$$

Its proof uses C3 and Central Limit Theorem.

Book: Alon & Spencer  
The Probabilistic Method.

Distinct Sums.

A set of positive integers  $\{x_1, \dots, x_n\}$  has distinct sums, if  $\forall S \subseteq \{1, \dots, n\}$  the sums  $\sum_{i \in S} x_i$  are distinct. Let  $f(n)$  denote the maximum  $k$  for which  $\exists$  set

$\{x_1, \dots, x_k\} \subseteq \{1, \dots, n\}$  with distinct sums.

Example.  $\{1, 2, 2^2, 2^3, \dots, 2^{\log_2 n}\} \subseteq \{1, \dots, n\}$ .  $f(n) \geq \lfloor \log_2 n \rfloor + 1 \geq \log_2 n$ .

To obtain an upper bound for  $f(n)$ , first observe that

$$\sum_{i=1}^k x_i \leq \sum_{i=1}^k x_k \leq nk \text{ for } f(n)=k$$

This means that each sum is less than  $nk$ . Since all the sums are distinct, we get  $\forall S \subseteq \{1, \dots, k\}$ ,  $\sum_{i \in S} x_i \in \{1, 2, \dots, nk-1\}$

Hence

$$\left\{ \sum_{i \in S} x_i : S \subseteq \{1, \dots, k\} \right\} \subseteq \{1, 2, \dots, nk-1\}.$$

Thus, (since there are exactly  $2^k$  subsets of  $\{1, \dots, k\}$ )

$$\begin{aligned} 2^k &= |\left\{ \sum_{i \in S} x_i \mid S \subseteq \{1, \dots, k\} \right\}| \\ &\leq |\{1, 2, \dots, nk-1\}| < nk \end{aligned}$$

Therefore  $\frac{2^k}{k} = \frac{2^{f(n)}}{f(n)} < n \rightarrow k < \log_2 n + \log_2 \log_2 n + 2$ .

**Theorem.** Assume that the above settings hold. Then,

$$f(n) \leq \log_2 n + \frac{1}{2} \log_2 \log_2 n + 5$$

**Merten's Theorem.**  $\sum_{p \leq x} \frac{1}{p} \sim \frac{\log \log x}{\pi(x)} \quad \lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 1$ .

Example

$$\sum_{\substack{p \leq P \\ p \text{ known}}} \frac{1}{p} = 4.$$

**Proof (Erdős, Moser, '56).** Fix any set  $\{x_1, \dots, x_k\} \subseteq \{1, \dots, n\}$  with distinct sums.

Let  $\xi_1, \dots, \xi_k$  be independent and uniformly distributed r.v.s. on  $\{0, 1\}$ . Define

$$X = \sum_{i=1}^k \xi_i x_i - \text{r.v.}$$

Since  $\xi_i$ 's are uniformly dist., then  $\forall i \in \{1, \dots, k\}$ ,  $P\{\xi_i = 0\} = P\{\xi_i = 1\} = 1/2$ .

Then  $E(X) = \frac{x_1 + \dots + x_k}{2}$

$$\begin{aligned} \text{Var}(X) &= \sum_{i=1}^k \text{Var}(\xi_i x_i) = \sum_{i=1}^k x_i^2 \text{Var}(\xi_i) = \sum_{i=1}^k x_i^2 (E(\xi_i^2) - E(\xi_i)^2) \\ &= \sum_{i=1}^k x_i^2 \left( \sum_{x=0}^1 x^2 P(\xi_i = x) \cdot \frac{1}{2} \right) \end{aligned}$$

$$= \sum_{i=1}^k x_i^2 / 4 = \frac{x_1^2 + \dots + x_k^2}{4} \leq \frac{n^2 k}{4}$$

This gives  $\sqrt{\text{Var}(x)} \leq \frac{n\sqrt{k}}{2}$

By CI, we put

$$\Pr(|X - \mathbb{E}(X)| \geq \lambda \sqrt{\text{Var}(X)}) \leq 1/\lambda^2 \text{ for any } \lambda > 0. \text{ So,}$$

$$1 - \Pr(|X - \mathbb{E}(X)| < \lambda \sqrt{\text{Var}(X)}) \leq 1/\lambda^2 \Rightarrow \Pr(|X - \mathbb{E}(X)| < \lambda \sqrt{\text{Var}(X)}) \geq 1 - \frac{1}{\lambda^2}.$$

Observe that for any  $s \in \{1, \dots, n\}$

$$\Pr(X = s) \in \left\{0, \frac{1}{2^k}\right\}$$

Therefore,

$$\Pr(|X - \mathbb{E}(X)| \geq 2\sqrt{\text{Var}(X)}) = \Pr(X \in (\mathbb{E}(X) - 2\sqrt{\text{Var}(X)}, \mathbb{E}(X) + 2\sqrt{\text{Var}(X)}))$$

Hence,

$$\Pr(|X - \mathbb{E}(X)| \geq 2\sqrt{\text{Var}(X)}) \leq \frac{2\lambda \sqrt{\text{Var}(X)} + 2}{2^k}$$

Combining this with (\*), we get

$$1 - 1/2^k \leq \frac{2\lambda \sqrt{\text{Var}(X)} + 2}{2^k} \quad (\text{as } \sqrt{\text{Var}(X)} \leq \frac{n\sqrt{k}}{2})$$

$$\lambda \geq \left(\frac{\gamma^2 - 1}{\gamma^3}\right) \frac{2^k}{\sqrt{k}} - \frac{1}{2\sqrt{k}} \quad \text{Pick } \lambda = \sqrt{3}, \text{ the result follows.}$$

→ Chernoff's Bound

↳ Sidon Set

→ F.K.G Inequality ↳ additive N.T.