

### Problem Statement

You are given an  $m \times n$  grid. Each cell of the grid has a direction pointing to the next cell you should visit if you are currently in this cell. The direction in  $\text{grid}[i][j]$  can be:

- 1 which means go to the cell to the right. (i.e go from  $\text{grid}[i][j]$  to  $\text{grid}[i][j + 1]$ )
- 2 which means go to the cell to the left. (i.e go from  $\text{grid}[i][j]$  to  $\text{grid}[i][j - 1]$ )
- 3 which means go to the lower cell. (i.e go from  $\text{grid}[i][j]$  to  $\text{grid}[i + 1][j]$ )
- 4 which means go to the upper cell. (i.e go from  $\text{grid}[i][j]$  to  $\text{grid}[i - 1][j]$ )

Notice that there could be some signs on the cells of the grid that point outside the grid.

You will initially start at the upper left cell  $(0, 0)$ . A valid path in the grid is a path that starts from the upper left cell  $(0, 0)$  and ends at the bottom-right cell  $(m - 1, n - 1)$  following the signs on the grid. The valid path does not have to be the shortest.

You can modify the sign on a cell with cost = 1. You can modify the sign on a cell one time only.

Return the minimum cost to make the grid have at least one valid path.

### Input Format

The first line contains two integers  $n$  and  $m$ , where  $m$  is the columns of the mat matrix and  $n$  is the process of the mat matrix.

The second line contains the input of the matrix.

### Output Format

Return the minimum cost to make the grid have at least one valid path.

### Constraints

- $m == \text{grid.length}$
- $n == \text{grid}[i].length$
- $1 \leq m, n \leq 100$
- $1 \leq \text{grid}[i][j] \leq 4$

### Sample Testcase 0

#### Testcase Input

```
4 4
1 1 1 1
2 2 2 2
1 1 1 1
2 2 2 2
```

#### Testcase Output

```
3
```

#### Explanation

You will start at point  $(0, 0)$ .

The path to  $(3, 3)$  is as follows.  $(0, 0) \rightarrow (0, 1) \rightarrow (0, 2) \rightarrow (0, 3)$  change the arrow to down with cost = 1  $\rightarrow (1, 3) \rightarrow (1, 2) \rightarrow (1, 1) \rightarrow (1, 0)$  change the arrow to down with cost = 1  $\rightarrow (2, 0) \rightarrow (2, 1) \rightarrow (2, 2) \rightarrow (2, 3)$  change the arrow to down with cost = 1  $\rightarrow (3, 3)$

The total cost = 3.

### Sample Testcase 1

#### Testcase Input

```
2 2
1 2
2 4
```

#### Testcase Output

```
1
```

**CODE**

```

#include<bits/stdc++.h>
using namespace std;

vector<vector<int>> dirs={{0,1,1},{0,-1,2},{1,0,3},{-1,0,4}};

int minCost(vector<vector<int>>& grid) {
    int n=grid.size(),m=grid[0].size();
    vector<vector<int>> dp(n,vector<int>(m,99999));
    //source distance will be 0
    dp[0][0]=0;
    deque<pair<int,int>> dq;
    dq.push_back({0,0});
    while(!dq.empty()){
        pair<int,int> p=dq.front();
        dq.pop_front();
        //if we reach the destination
        if(p.first==n-1 && p.second==m-1) return dp[p.first][p.second];

        for(vector<int> dir:dirs){
            int x=dir[0]+p.first;
            int y=dir[1]+p.second;
            if(x<0 || y<0 || x>=n || y>=m) continue;
            int cost;
            //checking the direction indicated by grid[p.first][p.second]
            if(grid[p.first][p.second]==dir[2])
                cost=0;
            else
                cost=1;
            //if cost of adjacent cell is less insert it in the queue
            if(dp[x][y]>dp[p.first][p.second]+cost){
                dp[x][y]=cost+dp[p.first][p.second];
                //if cost is 0 insert it in front of deque
                if(cost==0)

```

```

34             dq.push_front({x,y});
35             //else insert it at back of deque
36             else
37                 dq.push_back({x,y});
38         }
39     }
40 }
41 return dp[n-1][m-1];
42 }
43 }
44
45 int main(){
46     int n,m;
47     cin>>n>>m;
48     vector<vector<int>> grid(n,vector<int>(m));
49     for(int i=0;i<n;i++){
50         for(int j=0;j<m;j++){
51             cin>>grid[i][j];
52         }
53     }
54     cout<<minCost(grid);
55     return 0;
56 }
```

## CODE EXPLANATION

```
#include<bits/stdc++.h>
using namespace std;

vector<vector<int>> dirs = {{0,1,1},{0,-1,2},{1,0,3},{-1,0,4}};
```

This part of the code includes the necessary C++ standard libraries and defines a 2D vector `dirs`. Each subvector within `dirs` represents a direction and its associated information. The three values in each subvector are as follows:

- The first value is the change in the row index (i.e., `delta_x`) when moving in this direction.

- The second value is the change in the column index (i.e., `delta_y`) when moving in this direction.
- The third value represents the direction itself, where 1 means go right, 2 means go left, 3 means go down, and 4 means go up.

```
int minCost(vector<vector<int>>& grid) {
    int n = grid.size(), m = grid[0].size();
    vector<vector<int>> dp(n, vector<int>(m, 99999));

    // Source distance will be 0
    dp[0][0] = 0;
    deque<pair<int, int>> dq;
    dq.push_back({0, 0});
```

- The `minCost` function takes a 2D vector `grid` as input, which represents the grid with directions.
- `n` and `m` are the number of rows and columns in the grid, respectively.
- `dp` is a 2D vector used for dynamic programming, initialized with a large value (99999) to represent infinity.
- The distance to the source cell (0, 0) is set to 0, indicating that it costs nothing to start from there.
- `deque<pair<int, int>> dq` is a double-ended queue (deque) of pairs of integers. It is used for BFS (Breadth-First Search) to explore the grid.

```
while (!dq.empty()) {
    pair<int, int> p = dq.front();
    dq.pop_front();

    // If we reach the destination
    if (p.first == n - 1 && p.second == m - 1)
        return dp[p.first][p.second];

    for (vector<int> dir : dirs) {
        int x = dir[0] + p.first;
        int y = dir[1] + p.second;

        if (x < 0 || y < 0 || x >= n || y >= m)
            continue;

        int cost;
```

```

// Checking the direction indicated by grid[p.first][p.second]
if (grid[p.first][p.second] == dir[2])
    cost = 0;
else
    cost = 1;

// If cost of adjacent cell is less, insert it in the queue
if (dp[x][y] > dp[p.first][p.second] + cost) {
    dp[x][y] = cost + dp[p.first][p.second];

    // If cost is 0, insert it in front of deque
    if (cost == 0)
        dq.push_front({x, y});
    // Else insert it at back of deque
    else
        dq.push_back({x, y});
}
}

return dp[n - 1][m - 1];
}

```

- This part of the code implements the BFS algorithm to find the minimum cost.
- The `while (!dq.empty())` loop continues as long as there are cells in the deque to explore.
- For each cell `(p.first, p.second)` extracted from the deque, the code checks if it has reached the destination `(n - 1, m - 1)`. If it has, it returns the minimum cost, indicating that a valid path has been found.
- The code then iterates through the `dirs` vector, which defines the four possible directions to move (right, left, down, up).
- For each direction, it calculates the new row index `x` and column index `y` based on the current cell `(p.first, p.second)`.
- It checks if the new indices `x` and `y` are within the valid bounds of the grid (not out of bounds).
- The `cost` variable is determined based on whether the current cell's direction matches the direction defined by `grid[p.first][p.second]`. If they match, the cost is set to 0; otherwise, the cost is set to 1, indicating that a direction change is needed.
- If the cost of reaching the adjacent cell `(x, y)` is less than the current cost stored in `dp[x][y]`, it updates the minimum cost in `dp[x][y]`.

- Depending on whether the cost is 0 (no direction change) or 1 (direction change), it adds the cell  $(x, y)$  to the front or back of the deque, respectively. This ensures that cells with cost 0 are explored first, helping find shorter paths more efficiently.
- Finally, the function returns the minimum cost to reach the destination cell  $(n - 1, m - 1)$ .

```
int main() {
    int n, m;
    cin >> n >> m;
    vector<vector<int>> grid(n, vector<int>(m));

    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < m; ++j) {
            cin >> grid[i][j];
        }
    }

    cout << minCost(grid);
    return 0;
}
```

- In the `main` function, the code reads the dimensions of the grid (`n` and `m`) and the grid elements from the input.
- It then calls the `minCost` function to calculate the minimum cost to make the grid have at least one valid path.
- Finally, it prints the result.

This C++ code correctly solves the problem of finding the minimum cost to ensure a valid path through the grid and efficiently handles direction changes and path exploration using BFS