Task: Given a unique set of data

```
{"srn": "555555",
"name": "Anastasia Rizzo",
"exercise1": {
    "p": "2685735181983467",
    "g": "2",
    "a": "3628323521",
    "b": "5915667893",
    "cipherText": {
    "encoded": "1580593562655238",
    "base64": "BZ2KndYaBg=="
    }
}
}
```

Question 4. Verify whether g is a generator for p. Provide a brief explanation and include the method from your code, as well as the runtime

```
Given the following:
```

```
Prime number p = 2685735181983467
Generator g = 2
```

Prime number **p** can be a special kind of prime called "safe" prime.

Number **p** is a safe prime if both **p** and (**p-1**) /2 are prime numbers. (Wiener, 2003)

If **p** is a safe prime, it will be much easier to verify if **g** is a generator for **p**.

To check if number **p** is a safe prime, the following formulas have been applied:

```
To find q: \mathbf{q} = (\mathbf{p}-\mathbf{1})/2

\mathbf{q} = (2685735181983467 - 1)/2 = 1342867590991733

Snippet from java code (presented in Appendix 4, page 13, 14):

\mathbf{q} = ((\mathbf{p}.\mathbf{subtract}(\mathsf{one})).\mathsf{divide}(\mathsf{two})); // \mathbf{q} = (\mathbf{p}-\mathbf{1})/2

To find d: \mathbf{d} = \mathbf{g}^{\mathbf{q}} \mod \mathbf{p}

\mathbf{d} = 2^{1342867590991733} \mod 2685735181983467 = 2685735181983466

Snippet from java code (presented in Appendix 4, page 13, 14):

\mathbf{d} = \mathbf{g}.\mathsf{modPow}(\mathbf{q},\mathbf{p}); // \mathsf{d} = \mathsf{g}^{\mathsf{q}} \mod \mathbf{p}

\mathbf{f} = 2^{2685735181983466} \mod 2685735181983467 = 1

Snippet from java code (presented in Appendix 4, page 13, 14):

\mathbf{f} = \mathsf{g}.\mathsf{modPow}(\mathsf{d},\mathbf{p}); // \mathsf{check} if \mathbf{g}^{\mathsf{q}} \mathsf{d} \mod \mathbf{p} is prime
```

Result: $\mathbf{q} = (\mathbf{p-1})/2$ is a prime number, therefore \mathbf{p} is a safe prime number and \mathbf{g} is a generator for \mathbf{p} .

Anastasia Rizzo

To calculate the **runtime**:

Snippet from java code (presented in Appendix 4, page 13, 14):

long current_local_time = System.currentTimeMillis(); //runtime
System.out.println("Runtime: " + (System.currentTimeMillis() - current_local_time) + "
millisecond/s"); // print runtime

Result: 2 millisecond/s

Question 5. Considering that *a* is Alice's private key and *b* is Bob's private key, compute their public keys and show how they can generate the same shared key. Include a brief explanation and the relevant code snippet.

Given the following:

Prime number **p** = 2685735181983467 Generator **g** = 2 Alice's private key **a** = 3628323521 Bob's private key **b** = 5915667893

To compute Alice's public key (x) and Bob's public key (y) the following formulas have been applied:

 $x = g^a \mod p$ and $y = g^b \mod p$, where "p" is a prime number, "g" is a generator, "a" is a private key.

```
So, that \mathbf{x} = 2^{3628323521} \mod 2685735181983467 = 1268048862489310
\mathbf{y} = 2^{5915667893} \mod 2685735181983467 = 1346574413247867
```

Snippet from java code (presented in Appendix 5, page 14, 15):

```
x = g.modPow(a, p); // Alice's public key
y = g.modPow(b, p); // Bob's public key y
```

Result:

Alice's public key **x** = 1268048862489310 Bob's public key **y** = 1346574413247867

To compute Alice's shared key (kA) and Bob's shared key (kB) the following formulas have been applied:

 $kA = y^a \mod p$ and $kB = x^b \mod p$, where "y" is Bob's public key, "x" is Alice's public key, "a" is a private key, "p" is a prime number.

```
So, that: kA = 1346574413247867 <sup>3628323521</sup> mod 2685735181983467 = 789708769021392 kB = 1268048862489310 <sup>5915667893</sup> mod 2685735181983467 = 789708769021392
```

Snippet from java code (presented in Appendix 5, page 14, 15):

```
kA = y.modPow(a, p); // Alice's shared key k
kB = x.modPow(b, p); // Bob's shared key k
```

Result:

Alice's shared key **k** = 789708769021392 Bob's shared key **k** = 789708769021392

Question 6. Decrypt the provided cipher text which has been encrypted with the shared key that you computed in Question 5. Include a brief explanation and the relevant code snippet.

Given the following:

Prime number $\mathbf{p} = 2685735181983467$

Generator $\mathbf{g} = 2$

Alice's private key **a** = 3628323521

Bob's private key **b** = 5915667893

Alice's public key x = 1268048862489310

Bob's public key **y** = 1346574413247867

Alice's shared key **k** = 789708769021392

Bob's shared key **k** = 789708769021392

Cypher text "encoded" c = 1580593562655238

To decrypt the provided cypherText "encoded" the following formulas have been applied: $xInv = k \mod p$, where "k" is a shared key, "p" is a prime number.

m = (c * k) mod p, where "c" is a provided cypher text, "k" is a shared key, "p" is a prime number.

So that:

xInv = 789708769021392 mod 2685735181983467 = 1139504415921727 m = (1580593562655238 * 789708769021392) mod 2685735181983467 = 495890948466

Snippet from java code (presented in Appendix 6, page 14, 15):

xInv = kA.modInverse(p); // Inverse of x in decryption part m = c.multiply(xInv).mod(p); // plain text "encoded"

Result:

m = 495890948466

To decode plainText "encoded" from BigInteger to plainText "base64" this snippet of code has been applied:

```
Base64.Encoder encoder = Base64.getEncoder();
byte[] bigIntegerBytes = BigInteger.valueOf(495890948466L).toByteArray();
String base64EncodedBigIntegerBytes = encoder.encodeToString(bigIntegerBytes);
System.out.println("Plain text base64: "+ base64EncodedBigIntegerBytes);
```

To decode from plainText "base64" to plainText "text" this snippet of code has been applied:

```
Decoder decoder1 = Base64.getDecoder();
byte[] bytes = decoder1.decode(base64EncodedBigIntegerBytes);
String decodedString = new String(bytes, UTF_8);
```

Result:

The plainText "text" is: sugar

Question 7. Suppose Alice and Bob want to generate a new set of keys. They decide that they should use a 17-digit prime instead. How would they go on about generating a new p and a corresponding generator g? Provide a brief explanation and include the relevant code snippet, as well as its runtime.

Need to generate a new set of keys:

17-digit prime p and a corresponding generator g.

First, a new random **17-digit prime p** was generated.

Second, a new random corresponding generator g was generated.

Third, if number p is a safe prime was checked.

Fourth, whether **g** is a generator for **p** was verified.

Finally, **runtime** was calculated.

To generate new **17-digit prime p**:

A random number generator has been applied.

Snippet from java code (presented in Appendix 7, page 16, 17):

```
Random r = new Random(); // random number r
int numDigits = 17; // number of digits(17)
double LOG_2 = Math.log(10)/Math.log(2);
int numBits = (int) (numDigits * LOG_2);
p = new BigInteger (numBits, r); // prime number p
```

Result: new **17-digit prime p** = 7216142699972837

To generate a corresponding **generator g**:

Random number generator has been applied.

Snippet from java code (presented in Appendix 7, page 16, 17):

```
int low = 3; // low range
int high = 10; // high range
int e = r.nextInt((high - low) + 1) + low; // calculations of int e value in range between 3 and
10; int e = Integer i = BigInteger g
```

Result: a corresponding generator g = 3

To check if number **p** is a safe prime the following formulas have been applied:

```
To find q: \mathbf{q} = (\mathbf{p}-\mathbf{1})/2
\mathbf{q} = (7216142699972837 - 1)/2 = 3608071349986418
Snippet from java code (presented in Appendix 7, page 16, 17):
\mathbf{q} = ((\mathbf{p}.\mathbf{subtract}(\mathbf{one})).\mathbf{divide}(\mathbf{two})); // \mathbf{q} = (\mathbf{p}-\mathbf{1})/2
To find d: \mathbf{d} = \mathbf{g}^{\mathbf{q}} \mod \mathbf{p}
\mathbf{d} = 3^{3608071349986418} \mod 7216142699972837 = 7216142699972836
Snippet from java code (presented in Appendix 7, page 16, 17):
\mathbf{d} = \mathbf{g}.\mathbf{modPow}(\mathbf{q},\mathbf{p}); // \mathbf{d} = \mathbf{g}^{\mathbf{q}} \mod \mathbf{p}
\mathbf{f} = 3^{7216142699972836} \mod 7216142699972837 = 1
Snippet from java code (presented in Appendix 7, page 16, 17):
\mathbf{f} = \mathbf{g}.\mathbf{modPow}(\mathbf{d},\mathbf{p}); // \mathbf{check} \text{ if } \mathbf{g}^{\mathbf{q}} \text{ mod } \mathbf{p} \text{ is prime}
```

Result: $\mathbf{q} = (\mathbf{p-1})/2$ is a prime number, therefore \mathbf{p} is a safe prime number and \mathbf{g} is a generator for \mathbf{p} .

To calculate the **runtime**:

```
Snippet from java code (presented in Appendix 7, page 16, 17):
long current_local_time = System.currentTimeMillis(); //runtime
System.out.println("Runtime: " + (System.currentTimeMillis() - current_local_time) + "
millisecond/s"); // print runtime
```

Result: 2 millisecond/s.

Question 8. Generate a new set of private and public keys for Alice and Bob, using the p and g you generated in Question 7. Encrypt your SRN with the shared key. Include a brief explanation and the relevant code snippet.

Note:

The Question 8 was split on 2 parts.

Part 1 – to generate private keys (random) (code in Appendix 8 part 1, page 17, 18);

Part 2 – to generate public keys; to encrypt my SRN number with the shared key (code in Appendix 8 part 2, page 19, 20).

Part 1: To generate a new set of private keys (random).

Need to generate:

a new set of **private keys** for Alice (a) and Bob (b):

First, new **6-digit random numbers a** and **b** was generated (in this example, 6 digits were used, but it this can be changed to any number of digits).

Second, **if** number **a** and **b** are primes, has been checked by using Time Complexity: O(VN). Third, **runtime** was calculated.

```
To generate new 6-digit random numbers a and b:
```

A random number generator has been applied.

Snippet from java code (presented in Appendix 8 part 1, page 17, 18):

```
Random r = new Random(); // random number r
int numDigits = 6; // number of digits(6)
double LOG_2 = Math.log(10)/Math.log(2);
int numBits = (int) (numDigits * LOG_2);
p = new BigInteger (numBits, r); // random number p
```

Result: new 6-digit number a = 300889 and new 6-digit number b = 288293.

To check **if** number **a** and **b** are primes, Time Complexity O(VN) was applied: Snippet from java code (presented in Appendix 8 part 1, page 17, 18):

```
}
// Time Complexity: O(vN)
static void OSQRTNMethod(long a){
  boolean isPrime = true; // boolean

// if a is not a prime
  for (int i = 2; i <= Math.sqrt(a); i++) {
     if(a%i==0) {
        System.out.println("Number " + a +" is not a prime ");
        isPrime = false;
        break;
     }
  }
  // if a is a prime
     if(isPrime)
        System.out.println("Number " + a +" is a prime ");</pre>
```

Result:

```
Alice's private key a = 300889 is prime
Bob's private key b = 288293 is prime
```

To calculate the runtime:

OSQRTNMethod(a);

```
Snippet from java code (presented in Appendix 8 part 1, page 17, 18):
```

```
long current_local_time = System.currentTimeMillis(); // runtime
System.out.println("Runtime: " + (System.currentTimeMillis() - current_local_time) + "
millisecond/s"); // print runtime
```

Result: 1 millisecond/s.

Part 2: To generate Alice's (x) and Bob (y) public keys; to encrypt my SRN number with the shared key.

Given the following:

Prime number **p** = 7216142699972837Generator **g** = 3

Alice's private key **a** = 300889 Bob's private key **b** = 288293

Original message **m** = 140359547 (*my SRN number*)

To compute Alice's public key (x) and Bob's public key (y) the following formula has been applied:

 $x = g^a \mod p$ and $y = g^b \mod p$, where "p" is a prime number, "g" is a generator, "a" is a private key.

```
So, that \mathbf{x} = 3^{300889} \mod 7216142699972837 = 6321511077026144
\mathbf{y} = 3^{288293} \mod 7216142699972837 = 2205468983963930
```

Snippet from java code (presented in Appendix 8 part 2, page 19, 20):

```
x = g.modPow(a, p); // Alice's public key x calculations
y = g.modPow(b, p); // Bob's public key y calculations
```

To compute Alice's shared key (kA) and Bob's shared key (kB) the following formula has been applied:

 $kA = y^a \mod p$ and $kB = x^b \mod p$, where "y" is Bob's public key, "x" is Alice's public key, "a" is a private key, "p" is a prime number.

```
So, that: \mathbf{kA} = 6321511077026144 ^{300889} mod 7216142699972837 = 5704511815613684 \mathbf{kB} = 2205468983963930 ^{288293} mod 7216142699972837 = 5704511815613684
```

Snippet from java code (presented in Appendix 8 part 2, page 19, 20):

```
kA = y.modPow(a, p); // Alice's shared key k calculations
kB = x.modPow(b, p); // Bob's shared key k calculations
```

To encrypt my SRN number with the shared key the following formulas have been applied:

```
\mathbf{r} = \mathbf{g}^{kA} \mod \mathbf{p}, where "g" is a generator, "kA" is a shared key, "p" is a prime number. \mathbf{z} = \mathbf{x}^{kA} \mod \mathbf{p}, where "x" is a public key, "kA" is a shared key, "p" is a prime number. \mathbf{c} = (\mathbf{m}_i * \mathbf{x}) \mod \mathbf{p}, where "m<sub>i</sub>" is a message, "x" is a public key, "p" is a prime number.
```

```
So, that:
```

```
\mathbf{r} = 3^{5704511815613684} \mod 7216142699972837 = 654424626614820
```

```
z = 6321511077026144 ^{5704511815613684} mod 7216142699972837 = 5566298450981226 c = (140359547 * 6321511077026144) mod 7216142699972837 = 3859422092333348
```

Snippet from java code (presented in Appendix 8 part 2, page 19, 20):

r = g.modPow(kA, p);// $r = g^kA mod p$ calculations; for further calculations, Alice's shared key "kA" will be used as a shared key.

```
z = x.modPow(kA, p); // z=x^kA mod p calculations
```

```
c = m.multiply(z).mod(p); // c=(m*z)mod p; cypher text calculations
```

To decrypt the provided cypherText "encoded" the following formulas have been applied:

 $xd = r^a \mod p$, where "r" is a cypher text pair, "a" is an Alice's private key, "p" is a prime number.

xInverse = xd mod p

md = (xInverse * c) mod p, where "xInverse" is an inverse of x, "c" is a provided cypher text, "p" is a prime number.

So that:

```
xd = 654424626614820 <sup>300889</sup> mod 7216142699972837 = 5566298450981226

xInverse = 5566298450981226 mod 7216142699972837 = 6928872166590396

md = (1580593562655238 * 789708769021392) mod 2685735181983467 = 140359547
```

Snippet from java code (presented in Appendix 8 part 2, page 19, 20):

```
xd = r.modPow(a, p); // xd=r^a mod p calculations

xInverse = xd.modInverse(p); // inverse of x=x mod p=x^-1 calculations

md = xInverse.multiply(c).mod(p); // md=(x^-1 * c) mod p calculations
```

To encrypt and decrypt cypher text with Base64:

Snippet from java code (presented in Appendix 8 part 2, page 19, 20):

byte[] bytes = decoder1.decode(base64EncodedBigIntegerBytes);

// System.out.println("Plain text base64: "+ decodedString);

```
// Base64
```

```
Base64.Encoder encoder = Base64.getEncoder();
byte[] bigIntegerBytes = BigInteger.valueOf(2261795017164508L).toByteArray();
String base64EncodedBigIntegerBytes = encoder.encodeToString(bigIntegerBytes);
System.out.println("Plain text base64: "+ base64EncodedBigIntegerBytes);
Decoder decoder1 = Base64.getDecoder();
```

Result: CAkXMI2G3A==

String decodedString = new String(bytes, UTF 8);

Anastasia Rizzo

References

Irvin, J., 2019. BrutePrime.java. [Online]

Available at: https://gist.github.com/jonathan-irvin/22f8339849d0e5b44b6a

[Accessed 03 April 2019].

Wiener, J. M., 2003. Safe Prime Generation with a Combined Sieve, s.l.: s.n.