

## ADVANCED DATA STRUCTURE ASSIGNMENT

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APPLICATION

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- 1) A program P reads in 500 integers in the range [0..100] representing the scores of 500 students. It then prints the frequency of each score above 50. What would be the best way for P to store the frequencies?

Answer:

- Target Scores: Only scores between 51 and 100 are of interest, meaning we need to track a total of 50 scores.
- Efficiency: Using an array is optimal here as it allows quick access and updates, which is faster than other data structures like linked lists or hashmaps.
- Memory Use: A simple array of size 50 is sufficient to store the frequency of scores from 51 to 100.

**Steps to Implement:**

- Create a Frequency Array: An array of size 50 will be used, initialized with zero to count occurrences of scores between 51 and 100.

**int frequency[50] = {0};**

- Read and Process the Scores: Iterate over the 500 student scores and for each score above 50, update the corresponding index in the frequency array.
- Map Scores to Array Indices: For example, a score of 51 will be mapped to index 0 in the array, and a score of 100 will correspond to index 49.
- Print the Frequencies: After updating the array, print the frequencies for the scores that occurred at least once.

```
#include <stdio.h>
```

```
int main() {
    int frequency[50] = {0};
    int scores[500];
    // Example: Initialize scores array (replace with input logic)
    for (int i = 0; i < 500; i++) {
        scores[i] = rand() % 101;
    }

    // Process each score
    for (int i = 0; i < 500; i++) {
        if (scores[i] > 50) {
```

```

        int index = scores[i] - 51;
        frequency[index]++;
    }
}

// Print frequencies of scores above 50
for (int i = 0; i < 50; i++) {
    if (frequency[i] > 0) {
        printf("Score %d: %d occurrences\n", i + 51, frequency[i]);
    }
}

return 0;
}

```

### Explanation:

Step 1: The frequency array stores counts for scores between 51 and 100.

Step 2: Each score is processed, and if it is above 50, its corresponding index in the array is updated.

Step 3: The mapping is done by subtracting 51 from the score to determine its position in the array.

Step 4: The final frequencies are printed for scores that appeared at least once.

### Advantages of Using an Array:

- Fast Access: Updates and accesses are done in constant time.
- Memory Efficiency: Only 50 integers are needed to store the frequency counts.

2) Consider a standard Circular Queue 'q' implementation (which has the same condition for Queue Full and Queue Empty) whose size is 11 and the elements of the queue are q[0], q[1], q[2], ..., q[10]. The front and rear pointers are initialized to point at q[2]. In which position will the ninth element be added?

Answer:

In this problem, a circular queue of size 11 is given, with the elements stored in positions q[0] to q[10]. Both the front and rear pointers are initialized at q[2]. The goal is to determine the position of the ninth element to be added.

In a circular queue, both addition and removal of elements wrap around the array when the end is reached. The positions are handled using the front and rear pointers.

Modulo Operation: The rear pointer wraps around when it reaches the last position ( $q[10]$ ). This can be calculated using the modulo operation  $(\text{rear} + 1) \% \text{size}$ .

### **Step-by-Step Calculation:**

Initial Position: Both the front and rear pointers are at  $q[2]$ .

Element Addition: The first element is added at  $q[2]$  (the initial rear position), and after each addition, the rear pointer is updated as  $(\text{rear} + 1) \% 11$ .

Adding Nine Elements: The positions are calculated as follows:

1st element at  $q[2]$

2nd element at  $q[3]$

3rd element at  $q[4]$

4th element at  $q[5]$

5th element at  $q[6]$

6th element at  $q[7]$

7th element at  $q[8]$

8th element at  $q[9]$

9th element at  $q[10]$

Thus, the ninth element will be added at position  $q[10]$ .

3) Write a C Program to implement Red Black Tree

Answer:

A Red-Black Tree is a balanced binary search tree with additional properties that ensure the tree remains balanced, allowing for efficient insertion, deletion, and search operations. Each node in a Red-Black Tree contains an extra bit that signifies the color of the node, either red or black.

### **Key Properties of Red-Black Tree:**

Every node is either red or black.

The root is always black.

All leaves (NIL nodes) are black.

If a node is red, its children must be black (no two consecutive red nodes).

Every path from a node to its descendant leaves must have the same number of black nodes.

C Code Implementation:

```
#include <stdio.h>

#include <stdlib.h>

typedef enum { RED, BLACK } Color; typedef struct RBNode {
int data;
Color color;
struct RBNode *left; struct RBNode *right; struct RBNode *parent;
} RBNode;

RBNode* createNode(int data) {
RBNode* node = (RBNode*)malloc(sizeof(RBNode)); node->data = data;
node->color = RED; node->left = NULL; node->right = NULL; node->parent =
NULL; return node;
}

void leftRotate(RBNode **root, RBNode *x) { RBNode *y = x->right;
x->right = y->left;
```

```

if (y->left != NULL)
y->left->parent = x; y->parent = x->parent;
if (x->parent == NULL)
*root = y;
else if (x == x->parent->left) x->parent->left = y;
else
x->parent->right = y;
y->left = x;
x->parent = y;
}

void rightRotate(RBNode **root, RBNode *y) { RBNode *x = y->left;
y->left = x->right;
if (x->right != NULL)
x->right->parent = y; x->parent = y->parent;
if (y->parent == NULL)
*root = x;
else if (y == y->parent->right) y->parent->right = x;
else
y->parent->left = x;
x->right = y;
y->parent = x;
}

void fixViolation(RBNode **root, RBNode *z) { while (z != *root &&
z->parent->color == RED) {
RBNode *grandparent = z->parent->parent;
if (z->parent == grandparent->left) { RBNode *uncle = grandparent->right;
if (uncle != NULL && uncle->color == RED) { z->parent->color = BLACK;
uncle->color = BLACK; grandparent->color = RED; z = grandparent;
}

```

```

else {
    if (z == z->parent->right) { z = z->parent; leftRotate(root, z);
    }

    z->parent->color = BLACK; grandparent->color = RED; rightRotate(root,
    grandparent);
    }
    }

    else {
        RBNode *uncle = grandparent->left;
        if (uncle != NULL && uncle->color == RED) { z->parent->color = BLACK;
        uncle->color = BLACK; grandparent->color = RED; z = grandparent;
        }
        else {
            if (z == z->parent->left) { z = z->parent; rightRotate(root, z);
            }

            z->parent->color = BLACK; grandparent->color = RED;
            leftRotate(root, grandparent);
            }
            }
            }

            (*root)->color = BLACK;
            }

            void insert(RBNode **root, int data) { RBNode *z = createNode(data); RBNode
            *y = NULL;

            RBNode *x = *root;

            while (x != NULL) { y = x;
            if (z->data < x->data) x = x->left;
            else
            x = x->right;

```

```

}
z->parent = y; if (y == NULL)
*root = z;
else if (z->data < y->data) y->left = z;
else
y->right = z; fixViolation(root, z);
}
void inorder(RBNode *root) { if (root == NULL)
return;
inorder(root->left);
printf("%d (%s) ", root->data, (root->color == RED) ? "RED" : "BLACK");
inorder(root->right);
}
int main() {
RBNode *root = NULL;
int elements[] = {10, 20, 30, 15, 25, 40, 50};
int n = sizeof(elements) / sizeof(elements[0]);
for (int i = 0; i < n; i++) { insert(&root, elements[i]);
}
printf("Inorder Traversal of the Red-Black Tree:\n"); inorder(root);
printf("\n");
return 0;
}

```