SUPPLEMENTARY MATERIAL

C PSEUDO-CODE OF MDS AND MBPM

We provide the pseudo-code for MDS and MBPM in Alg. 6 and Alg. 7, respectively.

```
1167
              Algorithm 6: MDS
1168
1169
                 Input: (1) G = (V, E): an input graph
                            (2) p: E \rightarrow [0, 1]: activation probabilities
                            (3) S \subseteq V: a seed set
                            (4) b: an edge-removal budget
                            (5) h: the number of propagation hops
1173
                 Output: \mathcal{E} \subset E: a set of edges chosen to be removed
1174
              1 \text{ prob}(v) \leftarrow \pi(v; G, p, S)
                                                                                             ▶ MC simulation
1175
             _2 A \leftarrow the adjacent matrix of G
1176
              s rsa \leftarrow \sum_{0 \le i \le h} A^i \mathbf{1}
1177
             4 for i = 1, 2, ..., b do
1178
                       s_e \leftarrow (1 - \operatorname{prob}(u)) \cdot (\operatorname{rsa}(v) + \operatorname{rsa}(u)), \forall (v, u) \in E
1179
                       e \leftarrow \arg \max_{e \in E} s_e
1180
                       rsa \leftarrow update\_rsa(G, p, e, rsa)
1181
                       prob \leftarrow update\_prob(G, p, e, prob)
1182
                       E \leftarrow E \setminus \{e\}; \mathcal{E} \leftarrow \mathcal{E} \cup \{e\}
            10 return &
            11 Function update_prob(G, p, e = (s, t), prob):
1185
                       \Delta \text{prob}(v) \leftarrow 0, \forall v \in V
            12
1186
                       \Delta \texttt{prob}(t) \leftarrow \tfrac{p(s,t) \cdot \texttt{prob}(s) \cdot (1 - \texttt{prob}(t))}{1 - p(s,t) \cdot \texttt{prob}(s)}
            13
1187
                       U \leftarrow \{t\}; V \leftarrow \{t\}
            14
1188
                       while U \neq \emptyset do
            15
1189
                             U' \leftarrow \emptyset
             16
1190
                              foreach (u, v) \in E such that u \in U, v \in V do
             17
1191
                                    \Delta \mathsf{prob}(v) \leftarrow \Delta \mathsf{prob}(v) + \frac{p(u,v) \cdot \Delta \mathsf{prob}(u) \cdot (1 - \mathsf{prob}(v))}{1 - \mathsf{prob}(v)}
             18
1192
                                                                                    1-p(u,v)\cdot prob(u)
                                    if v \notin V then
             19
1193
                                     U' \leftarrow U' \cup \{v\}; V \leftarrow V \cup \{v\}
             20
1194
                             U \leftarrow U'
            21
1195
                       prob(v) \leftarrow prob(v) - \Delta prob(v), \forall v \in V
            22
                      return prob
            23
            Function update_rsa(G, p, e = (s, t), rsa):
1198
                       \Delta rsa(t) \leftarrow p(s,t) \cdot rsa(t)
            25
1199
                       U \leftarrow \{s\}; V \leftarrow \{s\}
            26
1200
                       while U \neq \emptyset do
            27
1201
                             U' \leftarrow \emptyset
            28
1202
                             foreach (v, u) \in E such that v \in V, u \in U do
             29
1203
                                    \Delta rsa(v) \leftarrow \Delta rsa(v) + p(v, u) \cdot \Delta rsa(u)
             30
                                    if v \notin V then
             31
1205
                                     U' \leftarrow U' \cup \{v\}; V \leftarrow V \cup \{v\}
             32
1206
             33
1207
                       rsa(v) \leftarrow rsa(v) - \Delta rsa(v), \forall v \in V
            34
1208
```

D ADDITIONAL EXPERIMENTAL RESULTS

D.1 Q1. Performances

return rsa

We report the effectiveness (the reduced ratio of influence) and the running time, along with the standard deviations, for each

Algorithm 7: Modified BPM

```
Input: (1) G = (V, E): an input graph
              (2) p: E \rightarrow [0, 1]: activation probabilities
              (3) S \subseteq V: a seed set
              (4) b: an edge-removal budget
              (5) d: the number of samplings
    Output: \mathcal{E} \subset E: a set of edges chosen to be removed
 0 \to 3
2 for i = 1, 2, ..., b do
         s_e \leftarrow 0, n_e \leftarrow 0 \forall e \in E
3
         for j = 1, 2, ..., b do
               l_e \sim \text{Bernoulli}(p(e)), \forall e \in E
 5
               E' \leftarrow \{e \in E : l_e = 1\}
               G' \leftarrow (V, E')
               n_S \leftarrow the number of nodes reachable from S on G'
               foreach e \in E such that l_e = 0 do
                s_e \leftarrow s_e + n_S; n_e \leftarrow n_e + 1
         foreach e \in E do
11
               if n_e \ge 1 then
12
                    s_e \leftarrow s_e/n_e
13
14
               else
               s_e \leftarrow \infty
15
         e \leftarrow \arg\min_e s_e
16
         E \leftarrow E \setminus \{e\}; \mathcal{E} \leftarrow \mathcal{E} \cup \{e\}
17
18 return {\cal E}
```

method, for each $b \in \{3, 5, 7, 10\}$, in Fig. 6 and Table 6. The standard deviations are computed over different seed sets and the values can be high since the seed sets can be highly different. Overall, it shows a fair trade-off between time and the reduction ratio except for RumorGuard-I.

D.2 Time for training and running GNN

We report the training time for a GCN in Table 7 and compare the running time of GNN with MC simulation in Table 8. Overall, the running time used by a GCN is at least 100 times shorter than the time taken by the MC simulation.

D.3 Budget scalability

We conduct experiments with all the considered methods with the budget b increasing from 1 to 10. In Table 9, for each method and each dataset, we report the minimum value of b with which the method runs out of time (i.e., takes more than one hour on a single seed set), where "." indicates that the method does not run out of time even with b=10. The methods whose outputs do not depend on the seed set are not included. Among the variants of Rumorguard, only Rumorguard-G fails on the largest dataset WC. Among the baseline methods, MBPM has the lowest budget scalability, and Greedy does not scale well on large graphs.

Table 6: The effectiveness (the reduced ratio of influence) of each method with the standard deviations, with budget $b \in \{3, 5, 7, 10\}$. O.O.T denotes out-of-time, i.e., the method does not terminate within one hour on a single seed set in the corresponding setting.

method	b = 3			b = 5		
	WC	CL	ET	WC	CL	ET
RANDOM	0.0016 (0.0214)	-0.0133 (0.1105)	0.0010 (0.0182)	-0.0045 (0.0216)	0.0010 (0.1188)	0.0081 (0.0158))
OdC	0.0022 (0.0225)	0.0063 (0.1077)	0.0123 (0.0207)	0.0117 (0.0238)	0.0352 (0.1534)	0.0119 (0.0407)
BC	-0.0006 (0.0226)	0.0006 (0.0860)	-0.0011 (0.0208)	0.0019 (0.0220)	-0.0057 (0.0992)	-0.0028 (0.0206)
PR	-0.0017 (0.0175)	-0.0036 (0.1204)	-0.0002 (0.0187)	-0.0017 (0.0201)	-0.0010 (0.1077)	-0.0124 (0.0182)
BPM	0.0.T	0.0159 (0.1221)	0.0387 (0.0471)	O.O.T	-0.0021 (0.1074)	0.0477 (0.0527)
KED	-0.0006 (0.0174)	0.0109 (0.1162)	-0.0015 (0.0180)	0.0029 (0.0178)	0.0097 (0.1317)	-0.0029 (0.0556)
MDS	0.0098 (0.0254)	-0.0004 (0.1085)	0.0298 (0.0366)	0.0173 (0.0229)	0.0202 (0.1341)	0.0368 (0.0421)
MBPM-100	0.0258 (0.1166)	0.0228 (0.1402)	-0.0010 (0.0187)	0.0198 (0.1173)	0.0032 (0.1289)	0.0140 (0.0844)
MBPM-1000	0.1253 (0.2175)	0.0026 (0.1236)	0.1249 (0.2049)	O.O.T	0.0107 (0.0936)	0.1710 (0.2209)
MBPM-10000	O.O.T	-0.0020 (0.1371)	0.3900 (0.2020)	O.O.T	0.0024 (0.1364)	0.4414 (0.2072)
MBPM-100000	O.O.T	O.O.T	0.4461 (0.1771)	O.O.T	O.O.T	O.O.T
GREEDY-10	0.0059 (0.0481)	-0.0182 (0.0976)	0.0244 (0.0995)	0.0062 (0.0459)	-0.0052 (0.1040)	0.0236 (0.1011)
Greedy-100	O.O.T	-0.0052 (0.1144)	0.1560 (0.2363)	O.O.T	0.0052 (0.1374)	0.1635 (0.2356)
RIS-0.6	0.2280 (0.1720)	0.3813 (0.1941)	0.2937 (0.1968)	O.O.T	0.5045 (0.2064)	0.3641 (0.2102)
RIS-0.4	O.O.T	0.4428 (0.1960)	0.3202 (0.2078)	O.O.T	0.5227 (0.2136)	0.4156 (0.2163)
RIS-0.2	O.O.T	0.4456 (0.1893)	0.3688 (0.2100)	O.O.T	0.5591 (0.2217)	0.4513 (0.2212)
RumorGuard-G	0.3609 (0.1796)	0.5240 (0.1876)	0.4675 (0.1855)	0.4311 (0.1726)	0.6547 (0.1877)	0.5613 (0.1828)
RumorGuard-O	0.3474 (0.1814)	0.5137 (0.1889)	0.4485 (0.1766)	0.4081 (0.1760)	0.6404 (0.2074)	0.5389 (0.1861)
RumorGuard-I	0.3261 (0.1728)	0.4783 (0.2062)	0.3920 (0.1869)	0.3747 (0.1717)	0.6149 (0.1877)	0.4703 (0.1780)

method	b = 7			b = 10		
	WC	CL	ET	WC	CL	ET
RANDOM	-0.0011 (0.0187)	-0.0187 (0.1244)	0.0014 (0.0210)	0.0002 (0.0201)	-0.0072 (0.1150)	0.0005 (0.0182)
OdC	0.0275 (0.0330)	0.0106 (0.1120)	0.0315 (0.0439)	0.0327 (0.0338)	0.0263 (0.1294)	0.0448 (0.0494)
BC	0.0027 (0.0194)	0.0196 (0.1118)	0.0090 (0.0557)	0.0020 (0.0234)	0.0054 (0.1013)	0.0158 (0.0534)
PR	-0.0035 (0.0225)	0.0187 (0.1071)	-0.0020 (0.0203)	0.0006 (0.0239)	0.0257 (0.1377)	0.0002 (0.0218)
BPM	0.0.T	0.0016 (0.0980)	0.0882 (0.1145)	0.0.T	-0.0085 (0.1030)	0.1047 (0.1432)
KED	-0.0010 (0.0158)	0.0017 (0.1316)	0.0079 (0.0569)	-0.0012 (0.0197)	-0.0100 (0.1319)	0.0113 (0.0579)
MDS	0.0160 (0.0244)	0.0015 (0.1112)	0.0372 (0.0442)	0.0163 (0.0258)	0.0230 (0.1354)	0.0396 (0.0424)
MBPM-100	0.0239 (0.1188)	-0.0071 (0.1172)	0.0160 (0.0862)	0.0238 (0.1185)	0.0003 (0.1092)	0.0176 (0.0860)
MBPM-1000	O.O.T	0.0084 (0.1020)	0.1759 (0.2193)	O.O.T	-0.0043 (0.1561)	0.1978 (0.2225)
MBPM-10000	O.O.T	-0.0074 (0.1214)	0.4741 (0.2028)	O.O.T	0.0248 (0.1461)	0.5046 (0.2000)
MBPM-100000	O.O.T	O.O.T	O.O.T	O.O.T	O.O.T	O.O.T
GREEDY-10	0.0076 (0.0493)	-0.0152 (0.1241)	0.0224 (0.1024)	0.0136 (0.0654)	-0.0300 (0.1453)	0.0321 (0.1111)
Greedy-100	O.O.T	0.0030 (0.1299)	0.1809 (0.2340)	O.O.T	-0.0074 (0.1414)	0.1977 (0.2358)
RIS-0.6	O.O.T	0.5732 (0.2292)	0.4172 (0.2134)	O.O.T	0.7590 (0.2020)	0.6390 (0.1443)
RIS-0.4	O.O.T	0.6093 (0.2117)	0.4688 (0.2131)	O.O.T	0.8113 (0.1836)	0.6517 (0.1534)
RIS-0.2	O.O.T	O.O.T	0.5116 (0.2195)	O.O.T	O.O.T	O.O.T
RumorGuard-G	O.O.T	0.7455 (0.1871)	0.6171 (0.1757)	O.O.T	0.8124 (0.1870)	0.6780 (0.1625)
RumorGuard-O	0.4513 (0.1705)	0.7303 (0.1838)	0.5918 (0.1787)	0.4967 (0.1607)	0.8023 (0.1824)	0.6477 (0.1702)
RumorGuard-I	0.4127 (0.1654)	0.6663 (0.2454)	0.5268 (0.1679)	0.4548 (0.1543)	0.7493 (0.1843)	0.5674 (0.1594)

Table 7: Average training time (in seconds) for training each GCN model.

WC	CL	ET	
17,227	7,390	6,818	

Table 8: The average running time (in seconds) for computing total influence per seed set when using GCN and MC simulations. MC simulations use 10,000 samplings.

model	WC	CL	ET	
GCN	0.0082	0.0056	0.0056	
MC simulation	3.7757	0.8276	0.7817	

0.75

0.50

0.25

0.00

0.0

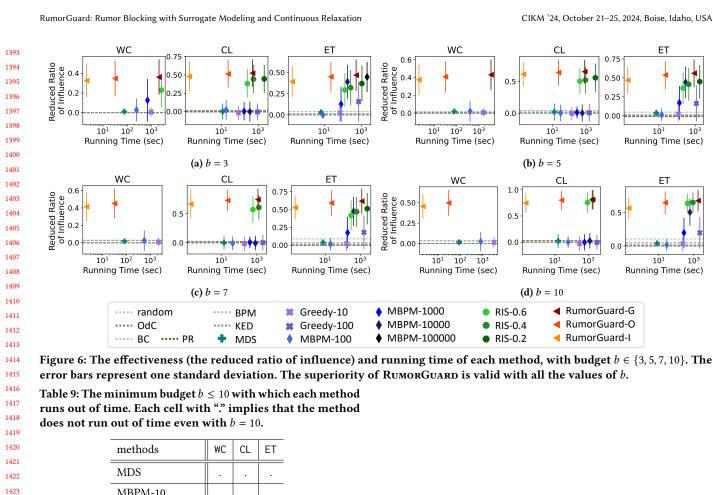
 10^{1}

Running Time (sec)

ΕT

Running Time (sec)

10³



error bars represent one standard deviation. The superiority of RUMORGUARD is valid with all the values of b.

WC	CL	ET
4		
1		
1	3	5
1		
5		
2		
1	6	10
6		