

Old programming idioms explained

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June 25, 2023

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1 Introduction

In the more than 60 years of its existence the Fortran programming language has undergone many changes, both in accordance with general insights in programming paradigms and in reaction to developments in computer hardware. This document discusses some of the idioms one finds in old FORTRAN packages and their modern alternatives.

Please be aware that the use of lowercase and uppercase forms of the name is not a whimsicality but rather marks a revolution within the language that was started with the publication of the Fortran 90 standard in 1990(?). Throughout this document we will use this distinction which is much more than merely typographic.

The set-up is simple and ad hoc: we discuss various idioms that have been used in the past decades and present contemporary equivalents or alternatives. Attempts are made to present them in a systematic way, but that mostly means grouping related topics.

2 Source forms

The source form of FORTRAN, now known as *fixed form*, shows its heritage in the era of punchcards: Each line could be up to 80 characters long, but only the first 72 characters had actual meaning. Some editors would add line numbers in the column 73–80 or people could add short comments. It had and has a few curiosities beyond the mere significance of the columns.

The alternative source form that was introduced with the Fortran 90 standard is more flexible and should definitely be used for any new source.

2.1 Fixed form versus free form

FORTRAN source files usually have an extension `.f` and on PCs, before you could use long names, the extension often was `.for`. There is, however, nothing special about these file extensions. They are merely a convention, useful for the compiler as it can use it to identify the source form. The free form source is often indicated by the extension `.f90`. *This should not be taken to mean that the code adheres to the Fortran 90 standard*, just as there is no particular difference as far as file extensions are concerned for FORTRAN 66 or FORTRAN 77.¹

To elaborate a bit on these file extensions:

- Fixed form is often identified by `.f` or for some old compilers on Windows `.for`, but that is not mandated by any standard and compilers often support options to specify what the source form is, if the extension is misleading.

¹File extensions themselves are a fairly recent feature of computer file systems. Older systems had different conventions.

- On file systems where file names are case-sensitive, such Linux, an extension `.F` or `.F90` is often meant to automatically invoke a preprocessor like C's preprocessor. Alternatively, you can also use compiler options to invoke the preprocessor explicitly.
- Preprocessing is not part of the FORTRAN or Fortran standards. It is simply an extension (no pun intended) that may or may not be supported by the compiler. As a consequence, there is no prescribed syntax, though very often the C conventions are used.

2.2 The significance of spaces

The original fixed form for FORTRAN sources has at least the following curiosities:

- In the columns 1 to 5 you can only put comments, if the first character is a comment character ("`C`" or "`*`"), or statement labels.
- The sixth column is reserved to indicate that the previous line is continued:

```
*234567890
      WRITE(*,*) 'A long sentence that spans over two',
&          'or more lines'
```

The continuation character, in the above the "`&`", can actually be any character with the exception of a zero (0). Some programmers would use digits, so that you can easily count the number of lines.

- The columns 7 to 72 were used for the actual code.
- The columns 73 to 80 were reserved for comments, to the discretion of the user.
- *Spaces had no significance*, except within literal strings.

Especially this last property may come as a surprise. To illustrate this (see `fixedform.f`):

```
PROGRAM FIXEDFORM

INTEGER :: I

DO 100 I = 1,10
    WRITE(*,*) I
100 CONTINUE
!
!   ILLUSTRATE THE FACT THAT SPACES HAVE NO MEANING ...
!
DO 110 I = 1,10
```

```

      I F ( I . E Q . 4 ) T H E N
        W R I T E ( * , * ) I
      E N D I F
110 C O N T I N U E
      E N D

```

Keywords in FORTRAN and Fortran are not reserved words and in fixed form you can put between the letters as many spaces you want, anywhere.

Running this program might produce the following output:

```

578772136
      4

```

”Might”, because there is an uninitialised variable here: the line `DO 100 I = 1.100` contains a typo. Instead of a comma, it contains a decimal point and thus the line is interpreted as:

```

DO100I = 1.10

```

The variable `DO100I` gets assigned the value 1.10, which is why there is no DO loop to produce ten lines in the output, counting from 1 to 10. Such mistakes can be caught by using the `IMPLICIT NONE` statement, a common enough extension.^[?]²

Note to self: <https://stevlionel.com/drfortran/2021/09/18/doctor-fortran-in-implicit-dissent/>

2.3 The character set

Officially, FORTRAN code should be written with capitals only. Lowercase letters are only allowed in literal strings. Again, a heritage of the machinery of old. But it has been allowed by compilers to use lowercase as well.

Beyond letters of the Latin alphabet you can use digits and underscores and several other characters. Even today, the standard is quite strict: characters outside the officially supported set are only allowed (or tolerated?) in literal strings.

2.4 Variables, functions, names and types

A feature that Fortran code relies on less and less is the use of *implicit typing*: of old a variable, or indeed a function, was of type **integer** if its name starts with one of: I, J, K, L, M or N. If not, the variable or function was of type **real**. You could declare the name as being of a different type, but that required an explicit declaration.

Because of this implicit typing, mistakes in names were a serious source of errors.

In FORTRAN, names were also limited to six (significant) characters, a limitation shared with the linkers.

²This statement was formalized in ”MIL-STD-1753”, precisely to make the coding safer.

This six-characters limit was even more severe as the names of subroutines, functions and `COMMON` blocks were global for the whole program. There was no other scope or control of visibility. This made it very important to explicitly define the names of such entities for any library you wrote, as naming conflicts could not be solved via a renaming clause, as they can nowadays in Fortran.

While the `IMPLICIT` statement is used nowadays mostly to turn off the implicit typing via `implicit none`, it can be used to define a particular type to variable whose names started with a particular character:

```
IMPLICIT DOUBLE PRECISION (A-H,O-Z)
```

means that all variables (and functions) whose names start with any of the letters A, B, ..., H or O, P, ... Z would have the type `DOUBLE PRECISION`, unless explicitly given a different type.

2.5 Special comments: D

Sometimes in old code you can see a character "D" in the first column. This is considered a comment line, unless you specify a compile option like `/d-lines` for the Intel Fortran compilers. Not all compilers support this, gfortran does not seem to have an option for it. The effect of specifying the relevant option is that the line is no longer considered a comment line, but is actual code. If supported by the compiler, it is only available for fixed form source.

It is a rather awkward form of "preprocessing":

- It is a compiler extension or at least not supported by all compilers. The now deprecated g77 compiler acknowledged the existence of such lines, but provided no facilities.

Note to self: <https://gcc.gnu.org/onlinedocs/gcc-3.4.6/g77/Debug-Line.html>

- It is not available for free form source.
- A practical problem is that it does not evolve with the code itself, as in the normal build process, these lines are considered comments.

It is just as easy to use ordinary code to provide debugging facilities:

```
*234567890
PROGRAM FIXEDDEBUG

D    WRITE(*,*) 'Note: in debugging mode ...'
    WRITE(*,*) 'Hello, world'

END PROGRAM FIXEDDEBUG
```

The program (see the file `fixeddebug.f`) is not accepted by gfortran because of the D-line, but Intel Fortran compiles it with and without the option `-d-lines`, only producing different output.

You could make the intent clearer using something along these lines (or a modern equivalent):

```

PROGRAM FIXEDDEBUG
LOGICAL DEBUG
PARAMETER (DEBUG = .FALSE.)

IF (DEBUG) WRITE(*,*) 'Note: in debugging mode ...'
WRITE(*,*) 'Hello, world'

END PROGRAM FIXEDDEBUG

```

With FORTRAN the `debug` parameter was best put in an include file to ensure that every program unit used the same setting. With Fortran a module containing such a parameter will do.

Note: If you use a *parameter* instead of a *variable*, many compilers will even eliminate the debugging code from the resulting program, but the code is compiled.

2.6 Compiler directives

Besides specific options, most if not all compilers have what are known as *compiler directives*. These take the form of a comment with a special format. As they are specific to the compiler, the precise form depends on the compiler, as well as what you can do with them. There are also compiler directives that are defined as a standard, like the *OpenMP* directives.³

Here is one example: an implicit way to redefine the default precision for real variables:

```

C For the Intel Fortran compiler:
!DIR$ REAL:8

```

This specifies the kind of real variables that are not declared with an explicit kind themselves.

Another example of a completely different nature:

```

C For the gfortran compiler:
!GCC$ unroll 10

```

The gfortran compiler supports loop unrolling and with this directive you control how this is done.

Note to self:

<https://gcc.gnu.org/onlinedocs/gfortran/UNROLL-directive.html>

<https://www.intel.com/content/www/us/en/docs/fortran-compiler/developer-guide-reference/2023-0/real-directive.html#GUID-38B61FE4-4058-497A-9AAA-2AF1BB56C178>

³Compiler directives have their equivalent in C. There they are called pragmas and function in much the same way.

Since these directives come with a specific compiler, you will need to check the documentation of the program or, more generally perhaps, the favoured compiler to see what the meaning is. It is not always entirely innocent: loop unrolling is an optimisation technique, but barring errors in the compiler, it should not matter for the outcome of a program whether the directive has been adhered to or not. This is different for the first directive – the default precision may have a significant effect!

2.7 Separate compilation

One of the objectives of the design of FORTRAN, which continues in today's standards, is that source files can be compiled independently. This should be read as: it is not necessary to compile the whole program in one step. In FORTRAN that was fairly easy: besides include statements there is no actual dependency possible. So, a source file can be compiled without knowing anything about other parts of the program.

In Fortran, there are dependencies beyond included files: if the program unit uses one or more (non-intrinsic) modules, then the order of compilation is that first the source files containing these modules must be compiled and then the source file with the using program unit.

The idea of separate compilation is a very powerful one, certainly in times when desk checking the code was a necessary first step, because the actual compilation might be done during the night, with you coming back in the morning to find out that you made a typo somewhere. Compiling source files separately and storing the resulting object files in libraries helps to minimize the risks of losing time.

But a side effect of this is that some programmers interpreted this as meaning: you should put all subroutines and functions in separate files. This does have one advantage, namely that the object files would only contain a single program unit and the linker would incorporate only the program units that are actually called. The disadvantage is that you may have to combine a large number of files into a library, because all the little subroutines and functions are in individual source files.

With the modules we have nowadays it is much easier to manage the routines. You can hide the ones that are not intended for outside use (so their names cannot cause trouble) and you can store the routines that functionally belong together in one and the same module, or organise the code in submodules.

3 Subroutines and functions

Older compilers were not so very sophisticated when it comes to checking the actual argument lists of calls to subroutines or functions against the dummy (expected) argument lists of these program units. This is partly due to there not being a mechanism in FORTRAN for providing such checks, even though some compilers could implement these checks at run-time. It was also used

sometimes to allow for flexibility. For instance: if you do not care if an array holds integers or reals, because you are simply writing them to some binary file, then one routine would suffice to take care of both types (perhaps even including double precision numbers).

A more common situation: allow one-dimensional arrays to be used as two-dimensional or pass scalars to array dummy arguments. As long as you remain with the allotted memory, there should be no problem.

3.1 Passing arrays

Several surprising idioms are connected to the passing of arrays to subroutines and functions:

- The dummy argument could be declared as `REAL ARRAY(1)`. This was not unusual in programs predating FORTRAN 77, as that introduced the asterisk for arrays whose sizes are not defined in the routine perse: `REAL ARRAY(*)`, the so-called assumed-size arrays. A small example:

```

SUBROUTINE PRINT( A, N )
  INTEGER N
C
C   Array declared pre FORTRAN 77:
C
  REAL A(1)
C
C   Array declared the FORTRAN 77 way:
C
C   REAL A(*)
  ...
END
```

In both cases, the routine would have to know what the actual size of the array should be. Often, that was done via an extra argument like the argument `N` in the example (see for a more elaborate and practical example section 6.6).

With Fortran you can use assumed-shape arrays, so that the dimensions of the array are passed on automatically. This does require that the interface is explicit. If the compiler does not see that interface, it has to assume that the FORTRAN 77 style is to be applied.

- The actual argument could be some element of the array:

```

REAL ARRAY(100)

CALL PRINT( ARRAY(10), 30 )
END
```

```

SUBROUTINE PRINT( A, N )
  INTEGER N
  REAL A(N)
  WRITE(*,*) A
END

```

The call to PRINT passes the array elements 10, 11, ... to the subroutine and in the subroutine it is assumed that the array A (its first dummy argument) is N (its second dummy argument) elements large. Then we can print the dummy array as a whole.

With Fortran we can specify a section of an array, not necessarily a contiguous section:

```

real :: array(100)

call print( array(10::10) )

```

The array section selects elements 10, 20, ... 100, but as the subroutine PRINT is in old-style FORTRAN, it does not get an *array descriptor*. The compiler will have to make a temporary array, copy in the selected elements, call PRINT and before releasing that temporary array, copy the elements of that temporary array into the selected elements. *It does not know the intent of the arguments of the subroutine PRINT after all!*

Marking the *intent* of arguments has at least two advantages: documenting the intended use of the argument and making certain optimisations possible. Like in the above example: copying back is not necessary if the argument is `intent(in)`.

3.2 Specific names for intrinsic functions

With the advent of FORTRAN 77 the language gained generic names of such intrinsic functions as MAX and SIN. Pre-FORTRAN 77 code uses specific names to distinguish the types of the arguments and the result. Thus:

- The minimum and maximum functions for integers were MIN0 and MAX0.
- The minimum and maximum functions for single-precision reals were AMIN1 and AMAX1.
- The minimum and maximum functions for double-precision reals were DMIN1 and DMAX1.
- If you wanted an *integer* result from a list of *reals*, then the function names were MIN1 and MAX1. But a *real* result from a list of *integers* is returned by AMIN0 and AMAX0.

- For the trigonometric functions: `SIN`, `DSIN` and `CSIN` would return the sine of respectively a single-precision real, a double-precision real and a complex value.

The Internet, as usual, will provide detailed information on the various intrinsic functions.

Sometimes, these specific names are required if you want to pass them as arguments to a routine. The best way is to avoid doing so and write a small helper function instead, as these specific names are deprecated and may even be removed from the standard altogether.

3.3 The `EXTERNAL` statement

Consider the following program:

```
PROGRAM USEFNC
  REAL A
  n = 10
  WRITE(*,*) A(N)

  END
```

You cannot build it, because it refers to a *function* `A` which is external to the program. You get an error about an unresolved symbol or a similar message. If you add the definition of a function `A` to the build step, then, of course, that symbol is resolved and the program can be built.

But the more important thing to notice is that the call to `A` is indistinguishable as it is from an array `A`. To make the distinction clear, you can use the `EXTERNAL` statement. This is also useful for arguments you pass that are actually procedures (see `integral.f`):

```
PROGRAM CALC
  REAL A, INTGRL
  EXTERNAL A, INTGRL

  XBEGIN = 0.0
  XEND    = 1.0
  WRITE(*,*) INTGRL(A, XBEGIN, XEND)
  END

  FUNCTION A(X)
  REAL X
  A = X ** 2
  END

  REAL FUNCTION INTGRL(A, XB, XE)
  REAL A, XB, XE
```

```

EXTERNAL A
REAL DX

DX      = (XE - XB) / 10.0
INTGRL = 0.0
DO 100 I = 1,10
    X      = XBEGIN + (I-1) * DX
    INTGRL = INTGRL + A(X)
100 CONTINUE

INTGRL = INTGRL * DX
END

```

Note that the declaration of INTGRL as a *real* function is required because the implicit typing rules would otherwise make it an integer, with very strange results.

3.4 Statement functions

Short functions like the function A in the previous section can also be written as so-called *statement functions*. These appear among the general declarations:

```

PROGRAM CALC
REAL INTGRL
C
C Use a statement function instead of a regular function
C
    A(X) = X ** 2

    XBEGIN = 0.0
    XEND   = 1.0

    INTGRL = 0.0

    DO 100 I = 1,10
        X      = XBEGIN + (I-1) * DX
        INTGRL = INTGRL + A(X)
    100 CONTINUE

    INTGRL = INTGRL * DX
    WRITE(*,*) INTGRL(A, XBEGIN, XEND)
END

```

As a statement function is not allowed as an actual argument, the integration is now moved into the program itself instead implemented in a separate function.

Statement functions can be useful to abbreviate expressions that appear at multiple locations in the code of a routine, but *internal procedures* in Fortran are much more useful and recognisable.

3.5 Constant values as actual arguments

A peculiar phenomenon was possible in the old days: you could change the value of a literal constant. That is: if you were sloppy, you could change the value of, say, 10 into 20:

```
PROGRAM CONST

CALL DOUBLE( 10 )
WRITE(*,*) 10
END

SUBROUTINE DOUBLE( X )
INTEGER X

X = 2 * X
END
```

Nowadays, compilers will not pass the literal constant itself but a copy or something similar, so that the above example will do no harm. But if the implementation is such that 10 is actually a memory address that is readable and writeable, then havoc is raised with, as a bonus, difficult to track bugs.

3.6 The ENTRY statement

FORTRAN offered the ENTRY statement to provide alternative entries into a subroutine or function. Such statement make it possible to call the procedure with a different argument list, for instance. A slightly contrived example:

```
REAL FUNCTION SETACC( X )
REAL X

REAL SUM
SAVE SUM

SUM      = X
SETACC = SUM
RETURN

ENTRY NXTACC( X )
REAL X

C
C      Accumulate to the variable SUM defined in SETACC
C
SUM      = SUM + X
NXTACC = SUM
END
```

With the function `SETACC` you initialise an "accumulator" variable and with `NXTACC` you add a new value to it and get the result. In a program you could use it like this (see `accum.f`):

```

PROGRAM ACCUM
REAL X, SUM

OPEN( 10, FILE = 'accum.inp' )

READ( 10, * ) X
SUM = SETACC( X )

100 CONTINUE
READ( 10, *, END = 110, ERR = 110 ) X
SUM = NXTACC( X )
GOTO 100

110 CONTINUE
WRITE(*,*) 'Sum: ', SUM
END

```

There are doubtless other uses for `ENTRY`, but the Fortran standard offers modules and module variables for this sort of things instead.

3.7 Alternative returns

A deprecated feature in Fortran, alternative returns provide a means to transfer the program control from a routine to its caller to another statement than the one following. This feature could be used for error handling (see `errors.f`):

```

PROGRAM ERRORS
REAL X

X = 2.0
CALL PRSQRT( X, *900 )

X = -1.0
CALL PRSQRT( X, *900 )
STOP

C
C   Handle any errors
C
900 CONTINUE
WRITE(*,*) 'Routine PRSQRT expects a positive argument'

END

```

```

SUBROUTINE PRSQRT( X, * )
REAL X

IF ( X .GE. 0.0 ) THEN
    WRITE(*,*) 'The square root of ', X, ' is ', SQRT(X)
ELSE
    RETURN 1
ENDIF
END

```

Note: because the feature is deprecated (obsolescent), compilers may not properly check that the actual call has a right number of statement labels in the interface. In most cases a simple return code works quite as well (the section 5.4 offers more on the subject).

4 Floating-point numbers

Nowadays, most computers use the IEEE format for representing floating-point numbers. The two main types you will encounter are single-precision reals, occupying four bytes of memory, and double-precision reals, occupying eight bytes. Fortran of old, including FORTRAN, favours single-precision – without any decoration a literal number like 1.23456789 is single-precision, whereas many other languages use double-precision reals by default.

Back in the days before the IEEE standard was widely adopted [?], reals were represented in many different ways and also the arithmetic operations we normally take for granted were not fully portable. This led to all manner of complications if you wanted to make your program portable, and that was and is certainly a goal of Fortran.

4.1 REAL*4 and REAL*8

One of the many extensions that were added by compiler vendors was the use of an asterisk (*) to indicate the precision:

```

*
*   Single precision real
*
*   real*4 x
*
*   Double precision real
*
*   real*8 y
*
*   Single precision complex
*
*   complex*8 z

```

The number indicates the number of bytes that a real would occupy. This has never been part of a FORTRAN or Fortran standard. The *kind* feature in Fortran is much more flexible, as it can capture aspects of the representation of floating-point numbers beyond mere storage.⁴

Note: Such notation is sometimes used for integers and logicals as well. Again the *kind* feature is much more useful than merely indicating the storage size.

Note: I have used a Convex computer in the distant past that actually had two different types of floating-point numbers that both were single-precision. One was structured according to the IEEE standard and the other was a native format. The difference for all intents and purposes was the interpretation of the exponent. The native format was said to be a bit faster, but they occupied the same storage, four bytes.

For the same bit pattern the *values* differed by a factor 4.

4.2 Literals in the source code

One thing to keep in mind: if a literal number occurs in the source code, it is interpreted as it appears, independent of the context. For Fortran this has been standardised: an expression on the right-hand side is evaluated independently of the left-hand side. More concretely:

```
double precision pi = 3.14159265358979323846264338327950288419716939937510
```

may look to specify π in some 50 decimals, but to the compiler it is merely a slightly bizarre way of expressing it in *single-precision*, so actually only six or seven significant decimals. To get *double-precision*, you need to add a *kind* or, as it was in FORTRAN, a "d" exponent (with some excess decimals removed):

```
double precision pi = 3.141592653589793238462643d0
```

You can see the difference if you run this program:

```
program diff_double_precision
  implicit none

  double precision, parameter :: pi_1 = 3.141592653589793238462643
  double precision, parameter :: pi_2 = 3.141592653589793238462643d0

  write(*,*) 'Difference: ', pi_1 - pi_2
end program diff_double_precision
```

which prints (you may expect slight differences in the last few decimals with different compilers):

⁴The current standard defines a general model for representing real numbers. This encompasses the IEEE formats, but is in fact more general.

Difference: 8.7422780126189537E-008

Some old FORTRAN compilers seem to have been less strict about the dichotomy between the left-hand side and the right-hand side and would indeed interpret such literal numbers as double-precision.

Another thing to keep in mind is that many compilers, both new and old, allow for compiler options that turn the *default* precision for a variable declared as **real** into *double precision*. If a program relies on this behaviour, then you need to carefully check the code.

4.3 Input in the absence of a decimal point

Disk storage nowadays is all but endless, but this luxury did not exist in the old days. This may have been the reason for a little known or used feature in input of real numbers: if a string representing a real number does not contain a decimal point, then the *input format* may insert it.

Here is an example, using internal I/O to make it self-contained (see also the file `input_no_point.f90`):

```
program show_insert_point
  implicit none

  real :: x
  character(len=10) :: string

  string = '1234'

  read( string, '(f4.0)' ) x
  write(*,*) x

  read( string, '(f4.2)' ) x
  write(*,*) x
end program show_insert_point
```

It produces:

```
1234.00000
12.3400002
```

With the format in the second read statement a decimal point is inserted!

It may have been useful in the past but it does suggest that to avoid surprises, you better not use input format with a prescribed number of decimals.

5 Control structures

FORTRAN came with a small number of control constructs and it was quite usual to construct other control flows via **IF** and **GOTO** statements. It inherited

some constructs from its predecessors that are very uncommon nowadays: the arithmetic (or three-way) **IF** and the computed **GOTO**, as well as the **ASSIGN** statement. This part of the document highlights these ancient idioms.

5.1 Ordinary and nested DO-loops

The ordinary **DO** loop in **FORTRAN** looks like this:

```
      DO 110 I = 1,10
          ... do something useful ...
110 CONTINUE
```

The statement label 110 indicates the end of the **DO** loop and anything in between is repeatedly executed. The Fortran equivalent is, unsurprisingly:

```
do i = 1,10
    ... do something useful ...
enddo
```

But there are a few more things to say about these **DO** loops. First of all, the statement label needs not appear with a **CONTINUE** statement. It could very well be put on the last executable statement:

```
      SUM = 0.0
      DO 110 I = 1,10
110      SUM = SUM + ARRAY(I)
```

It can even be used for multiple, nested, **DO** loops:

```
      SUM = 0.0
      DO 110 J = 1,10
          DO 110 I = 1,10
110      SUM = SUM + ARRAY(I,J)
```

To skip a part of the calculation, you can use a **GOTO** statement, where in Fortran you would use a **cycle** or **exit** statement:

```
*
* Sum the positive elements only and only if the sum
* remains smaller than 1.0
*
      SUM = 0.0
      DO 110 I = 1,10
          IF ( ARRAY(I) .LE. 0.0 ) GOTO 110
          IF ( SUM .GT. 1.0 ) GOTO 120
          SUM = SUM + ARRAY(I)
110 CONTINUE
120 CONTINUE
```

The example is a little contrived, so that you can see the use of the `GOTO` statement for both `cycle` and `exit`. The modern equivalent becomes:

```
!
! Sum the positive elements only and only if the sum
! remains smaller than 1.0
!
sum = 0.0
do i = 1,10
    if ( array(i) <= 0.0 ) cycle
    if ( sum > 1.0 ) exit
    sum = sum + array(i)
enddo
```

Note that sharing statement labels in a nested `DO` loop makes it difficult to see what a statement `GOTO endlabel` should mean: skip an iteration or skip the rest of the inner `DO` loop:

```
      SUM = 0.0
      DO 110 J = 1,10
      DO 110 I = 1,10
          IF ( SUM .GT. 1.0 ) GOTO 110
110      SUM = SUM + ARRAY(I,J)
```

In FORTRAN 66 (also known as FORTRAN IV) there was a significant difference with the `DO` loop you find in current Fortran: a `DO` loop would always be run at least once! This is due to the location of the check on the iteration condition, whether it is put at the start or at the end of the `DO` loop. Many compilers still provide an option to allow for the FORTRAN 66 semantics,[?] which includes this feature:⁵

```
* With the right compiler options, print this line once!
      DO 110 I = 1,0
          WRITE(*,*) 'FORTRAN 66: ', I
110 CONTINUE
      WRITE(*,*) 'Current value of i:', i
```

With FORTRAN 66 semantics the sample program (see `f66_loop.f90`) prints:

```
FORTRAN 66:           1
Current value of i:    2
```

With modern semantics it prints:

```
Current value of i:    1
```

This can result in subtle but nasty differences, if you are unaware of what was meant!

⁵I could not find such a flag for the gfortran compiler, but for Intel Fortran oneAPI it is `-f66`.

5.2 Simulating a DO-WHILE loop

There was no explicit DO WHILE construct in FORTRAN, at least not in the standard. Therefore you would need to simulate it using any of the following methods:

A DO loop with a large upper bound:

```
* Find the right line in a file
  DO 110 I = 1,10000000
    READ( 10, '(A)' ) LINE
    IF ( LINE(1:1) .NE. '*' ) THEN
      GOTO 120
    ENDIF
  110 CONTINUE
  120 CONTINUE
* Found the start of the information, proceed
  ...
```

A combination of statement labels and GOTO – check at the start:

```
* Find the right line in a file
  READ( 10, '(A)' ) LINE
  110 CONTINUE
  IF ( LINE(1:1) .NE. '*' ) GOTO 120
  READ( 10, '(A)' ) LINE
  GOTO 110
  120 CONTINUE

* Found the start of the information, proceed
  ...
```

(This example is a bit artificial to keep it in line with the other two, but similar constructs with different processing definitely occur in practice!)

A combination of statement labels and GOTO – check at the end:

```
* Find the right line in a file
  110 CONTINUE
  READ( 10, '(A)' ) LINE
  IF ( LINE(1:1) .EQ. '*' ) GOTO 110

* Found the start of the information, proceed
  ...
```

A modern equivalent would either use the DO WHILE loop or the unlimited DO loop:

```

!
! Find the right line in a file
!
read( 10, '(a)' ) line
do while (line(1:1) == '*' )
    read( 10, '(a)' ) line
enddo

!
! Found the start of the information, proceed
!
...

Or:

!
! Find the right line in a file
!
do
    read( 10, '(a)' ) line
    if (line(1:1) == '*' ) exit
enddo

!
! Found the start of the information, proceed
!
...

```

The precise location of the check on the condition depends on what the purpose is and whether you can actually check it at the start of the loop, as with a `DO WHILE`, or whether you require some preliminary calculation first. If you want to convert old-style source code, beware that the logic may sometimes have to be reverted, particularly if the condition comes at the end of the loop.

Note for self: <https://www.intel.com/content/www/us/en/docs/fortran-compiler/developer-guide-reference/2023-0/f66.html>

5.3 Three-way IF statements and computed GOTOs

Two types of statements that are quite alien to what you find in modern-day programming languages are the three-way or arithmetic `IF` statement and the computed `GOTO` statement. The latter could be used to simulate a `select case` construct, the first on the other hand was, in modern eyes, an unusual predecessor of the `IF ... ELSE ... ENDIF` block.

A *computed GOTO statement* takes a list of statement labels and a single integer expression:

```
GOTO (100, 200, 300) JMP
```

```

100 CONTINUE
    WRITE(*,*) 'Jump: 1'
    GOTO 400
200 CONTINUE
    WRITE(*,*) 'Jump: 2'
    GOTO 400
300 CONTINUE
    WRITE(*,*) 'Jump: 3'
400 CONTINUE
    ... the rest ...

```

Depending on the value of this expression (the value of `JMP` in the above example, the control would jump to the `Nth` label. If the value was zero or lower, the `GOTO` would not be executed and the program control would simply continue with the next statement. This is the case too with a value that is larger than the number of statement labels. (See as an illustration the source file `computed_goto.f90`)

Since there is nothing special about the statement labels the control would jump to, you had to make sure to jump somewhere else after the handling of each case. In the example that is done by jumping to label 400.

The `select case` construct of Fortran is better behaved, as you do not have to take care of jumping to the end yourself and it is possible to select the case via strings as well as integer values or even ranges.

There is nothing particularly magic about the *three-way IF statement*. But you need to know how it works:

```

    IF (IVALUE) 100, 200, 300
100 CONTINUE
    WRITE(*,*) 'Value is negative - ', IVALUE
    GOTO 400
200 CONTINUE
    WRITE(*,*) 'Value is zero - ', IVALUE
    GOTO 400
300 CONTINUE
    WRITE(*,*) 'Value is positive - ', IVALUE
400 CONTINUE
    ... the rest ...

```

The action, a jump to one of the three statement labels, to be taken depends on the *sign* of the integer expression. Often two of the statement labels would be the same, as two possibilities are more common than three. To see it in action, see the source file `three_way_if.f90`. Both statement types still exist in Fortran, or at least in the compilers, to support old-style programs.⁶

⁶The arithmetic IF statement was deleted from the Fortran 2018 standard, as gfortran will report. Intel Fortran will tell you this when you specify the standard as `f18 - "stand"`, defaulting to the Fortran 2018 standard.

5.4 Jumping to the end

The GOTO statement was and is also used to jump to a completely different part of the program unit for reporting error conditions:

```
      SUBROUTINE PRSQRT( X )

      IF ( X .LT. 0.0 ) GOTO 900

      WRITE(*,*) 'Square root of X = ', X, ' is ', SQRT(X)
      RETURN

900   CONTINUE
      WRITE(*,*) 'X should not be negative - ', X
      STOP
      END
```

Such statements can be gathered at the end of the program unit so as not to clutter the code that deals with normal processing. If you want to avoid the GOTO statement, then the modern BLOCK construct will help:

```
subroutine print_sqrt( x )
  real :: x

  normal: block
    if ( x < 0.0 ) exit normal

    write(*,*) 'square root of x = ', x, ' is ', sqrt(x)
    return
  end block normal
  !
  ! Error processing
  !
  errors: block
    write(*,*) 'x should not be negative - ', x
    stop
  end block errors

end subroutine print_sqrt
```

Note: the **errors** block is merely introduced to syntactically distinguish the error handling from the normal handling. It does not influence in this form the flow of control.

Note to self: <http://www.u.arizona.edu/~rubinson/copyright-violations/Go-To-Considered-Harmful.html>
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5.5 The ASSIGN statement

The uses of `GOTO` so far have all been static: the statement labels were fixed in the code. While the `GOTO` statement is frowned upon and it can certainly make the control flow difficult to follow when you do not use it in an orderly fashion, there is a possibility to use "dynamic" labels, so that the `GOTO` effectively jumps to a varying location. This is achieved by the `ASSIGN` statement. It is not often used, as in most cases better and especially clearer constructs are possible, even in FORTRAN, but here is one possible case:

Suppose you have to compute something complicated in a number of places in a program unit, based on a large number of variables, so that using a subroutine or a function is awkward, as it leads to a very long argument list. Nowadays you can easily use an internal routine, but this was not the case with FORTRAN. So, with the `ASSIGN` statement you could store a location to return to, jump ahead to the complicated piece of code, jump back when done, and remain in the same routine. Here is a simple example:

```
SUBROUTINE( ICASE )
  A = 1.0
  B = 2.0
  C = 3.0
  IF ( ICASE .EQ. 1 ) ASSIGN 100 TO JMP
  IF ( ICASE .EQ. 2 ) ASSIGN 200 TO JMP
  IF ( ICASE .EQ. 3 ) ASSIGN 300 TO JMP
  GOTO 900
```

* Case 1: use the result in F

```
100 CONTINUE
  WRITE(*,*) 'Case ', ICASE, 'value is ', F
  GOTO 400
```

* Case 2: use the result in G

```
200 CONTINUE
  C = 4.0
  WRITE(*,*) 'Case ', ICASE, 'value is ', G
  GOTO 400
```

* Case 3: use the result in H

```
300 CONTINUE
  WRITE(*,*) 'Case ', ICASE, 'value is ', H
  GOTO 400
```

* All done, continue

```
400 CONTINUE
  WRITE(*,*) 'Done'
  RETURN
```



```

*
  900 CONTINUE
* We can use the local variables directly
    F = A + B + C
    G = A + B * C
    H = A * B + C

* We are done, so return to the "caller"
    GOTO JMP
  END

```

It is a useless and contrived example, but it is only meant to illustrate how a "dynamic" jump can be constructed – the part after statement label 900 is actually independent of whatever happens above it. You can extend it with new cases, without having to worry about the computational part.

5.5.1 The DO-loop with a real index variable

A peculiar feature of FORTRAN 77 was the introduction of a DO-loop governed by a real variable. This means that the precise values of the variable might influence the number of iterations, for instance. A simple program to illustrate this (see `real_do.f90`):

```

PROGRAM REAL_DO
  IMPLICIT NONE

  INTEGER I, N
  REAL    R

  DO 120 I = 1,10
    N = 0
    DO 110 R = 0.0, 1.0*I, 0.1*I
      WRITE(*,*) R
      N = N + 1
    ENDDO
    WRITE(*,*) 'Number of iterations: ', N
  ENDDO
END

```

The output of the program with the Intel Fortran (leaving out the values of R):

```

Number of iterations:      11
Number of iterations:      11
Number of iterations:      10
Number of iterations:      11
Number of iterations:      11
Number of iterations:      10

```

```

Number of iterations:      11
Number of iterations:      11
Number of iterations:      10
Number of iterations:      11

```

So, some DO-loops get 11 iterations, what you would expect with exact arithmetic and others get only 10. If you run this program with the gfortran compiler the result is a uniform 11 iterations – apparently the loop is controlled in a different way! Apart from the unpredictability of the number of iterations with one compiler, you can also not rely on the portability of the program.

6 Memory management

In the decades leading up to the FORTRAN 77 standard, memory management was simple: declare what memory you need via statically sized arrays and that is it. There was no dynamic memory allocation, at least not in the FORTRAN language. It may surprise you, but even the concept of an operating system that took care of the computer was fairly new, as illustrated by a 1971 book by D.W. Barron, titled "Computer Operating Systems". Quoting from the book's jacket:

As the operating system is becoming an important part of the software complex accompanying a computer system. A large amount of knowledge about the subject now exists, mainly in the form of papers in computer journals. It is thus time for a book that coordinates what is known about operating systems.

There are, however, a few aspects of FORTRAN that make the story a bit more complicated: `COMMON` blocks, `EQUIVALENCE` and the `SAVE` statement. All three will be discussed here.

6.1 COMMON blocks

Variables, be they scalars or arrays, are normally passed via argument lists between program units (the main program, subroutines or functions). This is the immediately visible part. But you can also pass variables via `COMMON` blocks. These constitute a form of global *memory*, but not of global *variables*, as a `COMMON` block merely allocates memory and the mapping of memory locations onto variables is up to the program units themselves. For instance:

```

SUBROUTINE SUB1
COMMON /ABC/ X(10)
...
END

SUBROUTINE SUB2
COMMON /ABC/ A(5), B(5)

```

...
END

The `COMMON` block `/ABC/` appears in two subroutines, but in subroutine `SUB1` it is associated with the array `X` of 10 elements and in the subroutine `SUB2` it is associated with two arrays, `A` and `B`, both having five elements. The memory is shared, so that if you set `X(1)` to, say, 1.1 in subroutine `SUB1`, then on the next call to subroutine `SUB2`, the array element `A(1)` will have that same value, as they occupy the same memory location.

`COMMON` blocks should have the same *size* in all locations in the program's code where they occur. That is difficult to ensure, hence it was common (no pun intended) to put the declaration of `COMMON` blocks in so-called include files. Each program unit that needed to address the memory allocated via these `COMMON` blocks could then use the `INCLUDE` statement to have the compiler insert the literal text of that include file.⁷

There are various ways that `COMMON` blocks were used:

- Variables in `COMMON` blocks are persistent. At least, that was a very common occurrence. The rules in the FORTRAN standard are more complicated, but certainly with the `SAVE` statement you can rely on these variables to retain the values between calls to a routine.
- Often routines in a library have to cooperate: one routine is used to set options and other routines do the actual work. By using one or more `COMMON` blocks these options do not need to be passed around via the argument list.
- Together with `EQUIVALENCE` statements you could use the `COMMON` blocks to share workspace. Remember: back in the days memory was much and much more precious and scarcer than it is now. So, defining work arrays `WORK` (of type real) and `IWORK` (of type integer) and making them equivalent to each other, you could save on memory, if these arrays are not used at the same time.

In the code that would look like:

```
COMMON /ABC/ WORK(1000)
EQUIVALENCE (WORK(1), IWORK(1))
```

with a typical sloppiness with respect to array dimensions.

Nowadays, it is much easier to pass large amounts of essentially private data around, simply define a suitable derived type. Also, it is easy to allocate work arrays as you require them and release them again when done.

A special `COMMON` block was the so-called *blank* `COMMON` block. It had no name and it did not have to be declared with the same size in all parts of the program. In fact, on some systems it could be used as a flexible reservoir of memory, in much the same way as you have the heap nowadays. But this particular use was an extension to the standard.

⁷The `INCLUDE` statement was actually a common compiler extension.

6.2 More on EQUIVALENCE

TODO

6.3 The SAVE statement

According to the FORTRAN standard a local variable in a function or subroutine does not retain its value between calls, unless it has the **SAVE** attribute:

```
SUBROUTINE ACCUM( ADD )
*
*   Accumulate the counts
*
  INTEGER ADD

  INTEGER TOTAL
  SAVE      TOTAL
  DATA     TOTAL / 0 /

  TOTAL = TOTAL + ADD
  IF ( TOTAL > 100 ) THEN
    WRITE(*,*) 'Reached: ', TOTAL
  ENDIF
END
```

However, some implementations, notably on DOS/Windows, used static storage for these local variables, which meant that the variables would *seemingly* retain their values, even without the **SAVE** statement. If a program relied on this property and was ported to a different environment, all manner of havoc could be raised.

Note: I have actually had lively, but not necessarily pleasant, debates on whether the behaviour either way was correct. Sometimes the unexpected behaviour was claimed to indicate a compiler bug.

Some compilers have an option to enforce the **SAVE** attributes on variables, irrespective of the source code. You should take special care if old source code relies on such an option.

6.4 The initial values of (local) variables

A feature related in a way to the **SAVE** statement is the fact that in both FORTRAN and Fortran variables do not get a particular initial value, unless they have the **SAVE** attribute, implicitly or explicitly. With older compilers local variables may be stored in static memory and quite often they may have an initial value of zero or whatever the equivalent is for the variable's type, but that is in all cases simply a random circumstance. *Never assume that a variable that has not been explicitly given a value, has a particular value.*

You can set the initial value in FORTRAN via the **DATA** statement:

```

LOGICAL FIRST
DATA FIRST / .TRUE. /

```

This means that at the first call to the subroutine holding this variable `FIRST`, it has the value `.TRUE.`. You can later set it to `.FALSE.` to indicate that the subroutine has been called at least once before, so that no initialisation is needed anymore:

```

* Subroutine that sums the values we pass
SUBROUTINE SUM( x )
  INTEGER X

  INTEGER TOTAL
  LOGICAL FIRST
  DATA FIRST / .TRUE. /

  IF ( FIRST ) THEN
    FIRST = .FALSE.
    TOTAL = 0
  ENDIF

  TOTAL = TOTAL + X
END

```

(Just a variation on the previous example). This is not a very interesting routine, but it illustrates a typical use.

To emphasize: This type of initialisation is done so that the variables in question have the designated value at the first call. If you change the value, then they retain that new value. No reinitialisation occurs. (Actually, the value is not set on the first call, but rather is part of the data section of the program as a whole. There is no separate assignment.)

The `DATA` statement is not executable, it normally appears somewhere in the section that defines the variables, but it may occur elsewhere – most FORTRAN and Fortran compilers are not strict about it.

There is some peculiar syntax involved:

```

REAL X(100)
DATA X / 1.0, 98*0.0, 100.0 /

```

You can repeat values in much the same way as with edit descriptors in format statements: a count followed by an asterisk (*) and the value to be repeated. It is also possible to use implied do-loops:

```

INTEGER I
REAL X(100)

DATA (X(I), I = 1,100) / 100*1.0 /

```

While it is more usual to set the values together with the declaration of a variable nowadays, like:

```
integer :: i
real    :: x(100) = [ (1.0, i = 1,size(x)) ]
```

the DATA statement is more versatile, because it is not necessary to set the values for an array in one single statement:

```
integer :: i
real    :: x(100)

data (x(i), i = 1,50) / 50*1.0 /
data (x(i), i = 51,100) / 50*0.0 /
```

So, this old-fashioned statement may have its uses still.

Another peculiarity: the DATA statement has effect on the size of the object file and thus the executable itself. The following program leads to an executable of approximately 5.7 MB using gfortran on Windows:

```
program data_stmt
  implicit none

  integer :: i
  real    :: array(1000000)

  data (array(i), i = 1,size(array),2) / 500000*1.0 /
  data (array(i), i = 2,size(array),2) / 500000*2.0 /

  !
  ! Alternative
  !
  !! array(1::2) = 1.0
  !! array(2::2) = 2.0
  !

  write(*,*) sum(array)
end program data_stmt
```

If, instead, you use the alternative and remove the DATA statements, the executable is only 1.7 MB. Of course, this is an exaggerated example, but it illustrates that such DATA statements are very different in character than ordinary, executable statements.

6.5 Initialising variables in COMMON blocks: BLOCK DATA

The DATA statement plays an important role when it comes to initialising variables in a COMMON block. Since the COMMON blocks usually appear in more than

one subprogram (main program, subroutines, functions), they cannot be initialised in the same way as ordinary variables: which `DATA` statement should prevail, if several initialise the same `COMMON` variables?

Thus enter the `BLOCK DATA` program unit!

It is the only way to initialise variables in a `COMMON` block and it is special, because it is not executable and is not part of a routine or the main program. The peculiar consequence is that you cannot put it in a library: there is no reference to it, unlike with subroutines and functions, so it would never be loaded. Instead, you will normally put in the same file as the main program or link against its object file explicitly.

The general layout is:

```
BLOCK DATA
... COMMON blocks ...
... DATA statements ...
END
```

Here is a small example of this effect:⁸

```
gfortran -o common1 common.f90 block.f90
gfortran -o common2 common.f90
```

Both build commands succeed, despite the fact that one part of the program is missing in the second one.

The `block.f90` source file is:

```
BLOCK DATA
COMMON /ABC/ X
DATA X /42/
END
```

The `common.f90` source file is:

```
PROGRAM PRINTX
COMMON /ABC/ X
WRITE(*,*) 'Expected value of X = 42:'
WRITE(*,*) 'X = ', X
END
```

Program `common1` prints the value 42, whereas the other program prints 0. If the `BLOCK DATA` program unit had been an ordinary program unit, the building of this version would have failed on an unresolved symbol or the like.

⁸I use the `gfortran` compiler to illustrate such effects, but it would be similar with other compilers. And since most if not all FORTRAN features are still supported in Fortran, I also use free-form sources.

6.6 Work arrays

In the old days you could encounter arguments to a routine that represented such workspace. Usually you would have to declare the arrays to a size that matches the problem at hand. Here is an example from the *LAPACK* library for linear algebra:

```
      SUBROUTINE DGELS( TRANS, M, N, NRHS, A, LDA, B, LDB, WORK, LWORK,
$                   INFO )
*
*  -- LAPACK driver routine (version 3.2) --
*  -- LAPACK is a software package provided by Univ. of Tennessee, --
*  -- Univ. of California Berkeley, Univ. of Colorado Denver and NAG Ltd.--
*    November 2006
*
*  .. Scalar Arguments ..
      CHARACTER          TRANS
      INTEGER            INFO, LDA, LDB, LWORK, M, N, NRHS
*
*  .. Array Arguments ..
      DOUBLE PRECISION  A( LDA, * ), B( LDB, * ), WORK( * )
*
```

In this case, the argument `WORK` is a double-precision array of size `LWORK`. In the comments that document the use of this routine the precise usage is described:

```
*  WORK      (workspace/output) DOUBLE PRECISION array, dimension (MAX(1,LWORK))
*             On exit, if INFO = 0, WORK(1) returns the optimal LWORK.
*
*  LWORK      (input) INTEGER
*             The dimension of the array WORK.
*             LWORK >= max( 1, MN + max( MN, NRHS ) ).
*             For optimal performance,
*             LWORK >= max( 1, MN + max( MN, NRHS )*NB ).
*             where MN = min(M,N) and NB is the optimum block size.
*
*             If LWORK = -1, then a workspace query is assumed; the routine
*             only calculates the optimal size of the WORK array, returns
*             this value as the first entry of the WORK array, and no error
*             message related to LWORK is issued by XERBLA.
```

This means that for a particular case you can either use one of the formulae or the special value `-1` for `LWORK` to obtain an optimal value. The work array itself would still be a statically declared array.

Note that with the current features of Fortran the interface could be greatly simplified:⁹

⁹Intentionally left in fixed form.


```

SUBROUTINE DGELS( TRANS, A, B, INFO )
*
*   .. Scalar Arguments ..
CHARACTER, INTENT(IN) :: TRANS
INTEGER, INTENT(OUT) :: INFO
*
*   ..
*   .. Array Arguments ..
DOUBLE PRECISION, INTENT(INOUT) :: A(:, :), B(:, :)
*
*   ..

```

provided the interface is made explicit via a module or an interface block.

The careful reader will note that one feature of the original interface has not been retained in the simplification: `NHRS`. Thus, with this revised interface the array `B` should consist entirely of right-hand side vectors.

7 Input and output intricacies

Placeholder:

- command-line arguments for file names
- big-endian and little-endian
- list-directed input and output - also: /
- narrow formats (?)
- use of `d00` in input
- FORTRAN 66 semantics: `OPEN - STATUS = 'NEW'` as default.
- Effect of `BLANK = 'ZERO'` versus `BLANK = 'NULL'`
- 32-bits machines and unformatted files

Almost any program will read some kind of input files and produce some kind of output files. FORTRAN defined, roughly, three types of files:

- Text files meant to read or edited by humans. These are known as *formatted files*.
- Binary files with a record structure of sorts, so that you could read a part of the record and then automatically jump to the next. These are *unformatted, sequential files*. They are compact and might be considered *binary files*.
- Binary files with records that have a fixed length and where you can position the `READ` or `WRITE` action to a particular record: the *unformatted direct-access files*.¹⁰

¹⁰Actually, there are also formatted direct-access files, but these are seldom used.

Input and output in FORTRAN was always oriented towards records. For instance, if you read a number from a formatted file, then the read position is automatically move to the start of the next line, independent of the amount of data left on the previous line.

Similarly, writing to a file always produced a complete line. And the next write action would start on a new line.

Since Fortran 2003 the language also supports *stream-access* for both unformatted and formatted files. Before that standard, there were several more or less popular extensions to achieve the same effect.

7.1 Standard input and output

FORTRAN had no concept of standard input and standard output files. But there was a de facto standard in the form of logical unit numbers that were predefined:

- The unit number 6 was typically connected to the terminal:

```
WRITE(6,*) 'Enter a number:'
```

- The unit number 5 was typically connected to the keyboard:

```
*      Let the user type a number and read it
      READ(5,*) value
```

- Other "magic" unit numbers could exist, such as 7, often connected to a tape drive.
- Standard error could be connected to unit number 0 on some systems and would behave differently than standard output: no buffering, not redirected to file without extra care.

The FORTRAN 77 standard introduced the asterisk as a convenient way to define standard input and output without referring to a specific number: even though almost universal, these numbers were not defined by the FORTRAN standard.

With the advent of Fortran 2008 the intrinsic module `iso_fortran_env` now defines parameters like `INPUT_UNIT` to indicate the logical unit number for standard input and similar ones for standard output and error.

7.2 Carriage control

Printers of old had a simple system to control the positioning of the output on the paper: a character in the first column of a line to be output determined what would happen. Here are the (common) codes:

On some systems it would work on the terminal as well.

An extension to the write statement that had to do with controlling the cursor on a terminal was the dollar sign:

| Character | Meaning |
|-----------|--|
| space | Go to the next line – the usual way of positioning |
| 0 | Skip one line |
| 1 | Go to the next page |
| + | Do <i>not</i> go to the next line – print over the current |

```
WRITE(6,'(a,$)') 'Enter a number:'
```

The dollar sign meant that the cursor should stay on the last written line.

7.3 Direct-access files

By default, direct-access files are unformatted – values are dumped to the file or retrieved without the help of a human-readable format. Properties of direct-access files:

- You have to specify a record length, which holds for all records, when opening such a file. You can, however, open the file with different record lengths, if your application benefits from that. In other words, the length is not a property of the file itself.
- Direct-access files can be read or written by specifying a record number. Thus, like the name implies, you can jump around in the file at will.
- The record length is usually specified in *bytes* but some compilers, like Intel Fortran, use *words* as the unit. Where there was no way to know programmatically in FORTRAN what size was meant, since Fortran 2003, the intrinsic module `iso_fortran_env` contains the parameter `file_storage_unit` which is the size in bits.

Direct-access files, due to their simplicity, are compact and portable, as the structure of unformatted files depends on the compiler that was used for building the program (see below).

The main issues that makes these files non-portables are the binary representation of the numbers they contain. Nowadays, the main variation is the *endianness*: the order of the bytes that make up the number.

Here is a simple example of opening, writing and reading a direct-access file:

```
PROGRAM DIRECTACCESS

REAL VALUE(10)
INTEGER I, REC

OPEN( 10, FILE = 'directaccess.bin', ACCESS = 'DIRECT',
&      RECL = 4*10 )

*
```

```

* CALCULATE SOME DATA AND WRITE THEM TO THE FILE
*
      DO 120 REC = 1,10
        DO 110 I = 1,10
          VALUE(I) = I + REC * 10.0
110      CONTINUE

        WRITE( 10, REC = REC ) VALUE
120 CONTINUE

*
* READ THE DATA - IN REVERSE ORDER
*
      DO 220 REC = 10,1,-1
        READ( 10, REC = REC ) VALUE(1), VALUE(2)
        WRITE( *, * ) VALUE(1), VALUE(2)
220 CONTINUE
      END PROGRAM

```

It produces output like:

| | |
|------------|------------|
| 101.000000 | 102.000000 |
| 91.0000000 | 92.0000000 |
| 81.0000000 | 82.0000000 |
| 71.0000000 | 72.0000000 |
| 61.0000000 | 62.0000000 |
| 51.0000000 | 52.0000000 |
| 41.0000000 | 42.0000000 |
| 31.0000000 | 32.0000000 |
| 21.0000000 | 22.0000000 |
| 11.0000000 | 12.0000000 |

7.4 Unformatted sequential files

Data in sequential files, as the name suggests, are accessed in the order in which they appear in the files. For unformatted files you write the records one by one and you can read the records back one by one. But you cannot read more data from the record than its length. This is actually encoded in the file. The following program will therefore fail, as it tries to read more data than present in the first record:

```

PROGRAM SEQUNFORM

REAL VALUE(10)
INTEGER I, J

OPEN( 10, FILE = 'sequnform.bin', FORM = 'UNFORMATTED' )

```

```

*
* CALCULATE SOME DATA AND WRITE THEM TO THE FILE
* THE RECORDS GET LONGER
*
      DO 120 J = 1,10
        DO 110 I = 1,10
          VALUE(I) = I + J * 10.0
110      CONTINUE

          WRITE( 10 ) (VALUE(I), I = 1,J )
120 CONTINUE

*
* READ THE DATA - IN REVERSE ORDER
*
      REWIND( 10 )

*
* FIRST RECORD: ONLY ONE VALUE
*
      READ( 10 ) VALUE(1)
      WRITE( *, * ) VALUE(1)

*
* SECOND RECORD: TWO VALUES, BUT READ TEN
*
      READ( 10 ) VALUE
      WRITE( *, * ) VALUE(1)

      END PROGRAM

```

Reading the first record works, but it fails on the second, as the program tries to read 10 values, whereas the record only contain two (the gfortran was used):

```

11.0000000
At line 38 of file sequform.f (unit = 10, file = 'sequform.bin')
Fortran runtime error: I/O past end of record on unformatted file

```

Error termination. Backtrace:

```

Could not print backtrace: libbacktrace could not find executable to open
#0  0xa11301fa
#1  0xa11278a1
#2  0xa1122d90

```

```
#3  0xa1148e3e
#4  0xa1130e81
#5  0xa1101841
#6  0xa11018e4
#7  0xa11013bd
#8  0xa11014f5
#9  0xb56c7613
#10 0xb76a26a0
#11 0xffffffff
```

One essential difference between unformatted sequential files and stream-access files, which were introduced in the Fortran 2003 standard, is that these unformatted sequential files hold extra bytes to indicate the records. These extra bytes are an implementation detail, not defined in the standard. A common scheme is:

- The record starts with an integer number of four bytes that defines how many bytes will follow.
- The bytes after that integer are the actual data.
- The record is closed with the same integer number, so that a check is possible but it is also possible to go back in the file via the `BACKSPACE` statement¹¹

Thus the following statement has the effect of skipping a record:

```
READ(10)
```

whereas with a stream-access file nothing would happen, as these files have no concept of a record.

A common extension was that of the so-called *binary* file. Such a file is essentially what we now know as stream-access, but opening a file as *binary* required a specific, compiler-dependent keyword and not all compilers supported this type. Here is an example:

```
OPEN(10, FILE = 'data.bin', FORM = 'BINARY')
```

7.5 Reading an array

FORTTRAN 77 defined a convenient way to read in arrays:

```
REAL A(100)
...
READ(10,*) A
```

This reads 100 numbers from the file connected to unit number 10. But you can also use the following, using a so-called implied do-loop:

¹¹The `BACKSPACE` statement originated with tapes as main mass data storage.

```

INTEGER I, N
REAL A(100)
...
READ(10,*) N
READ(10,*) (A(I), I = 1,N)

```

A clever trick is to read the number of values and the values themselves in one statement:

```

INTEGER I, N
REAL A(100)
...
READ(10,*) N, (A(I), I = 1,N)

```

Because the value of N is known when the do-loop starts, there is no problem reading the array in this fashion. (Unless, of course, N exceeds the size of the array.)

8 Subjects

- array(*) versus array(:)
- array(10) as the starting point
- history of computers:
 - hardware
 - memory management
 - tools like source code control systems
 - connections between computers, Internet
- equivalence
- constants as actual arguments
- intent
- temporary arrays - non-contiguous arrays
- implicit types
- double precision versus kind
- checking interfaces
- separate compilations, the misunderstanding of one routine per file
- fixed form and spaces
- standard input and output
- LU-numbers 5 and 6 (and 7)
- command-line arguments for file names
- real do-variables
- entry
- statement functions
- six characters
- numerical binary representations versus IEEE (IBM, Cray, Convex)
- big-endian and little-endian
- double complex

- unformatted versus binary files
- list-directed input and output - also: /
- narrow formats (?)
- use of d00 in input
- Cray pointers
- uppercase/lowercase letters
- D as comment character
- external, also notation "*tan"
- FORTRAN 66 semantics: OPEN - STATUS = 'NEW' as default.
- Effect of BLANK = 'ZERO' versus BLANK = 'NULL'
- Specific names for functions like max() and sin()
- Alternate return
- Equating logicals: logical x, y; if (x .eq. y) then
- r1mach and friends
- Hollerith
- DECODE/ENCODE:

From discourse: (19 june 2023)

```
CHARACTER S*6 / '987654' /, T*6
INTEGER V(3)*4
DECODE( 6, '(3I2)', S ) V
WRITE( *, '(3I3)') V
ENCODE( 6, '(3I2)', T ) V(3), V(2), V(1)
PRINT *, T
END
```

The above program has this output:

```
98 76 54
547698
```

The DECODE reads the characters of S as 3 integers, and stores them into V(1), V(2), and V(3). The ENCODE statement writes the values V(3), V(2), and V(1) into T as characters; T then

References to be added: IEEE 754 and Fortran 90 standard.