



Workshop on the Integration of Information
Retrieval and Databases (WIRD'04)

An XML-IR-DB Sandwich



Is it better with an Algebra in
Between?

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Outline

- # Motivation
- # XML and Relational Databases
- # Region Algebra & Operator Properties
- # Region Algebra & Relevance Ranking
- # Properties of Ranking Operators
- # Conclusions and Future Work

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XML-IR and Relational DBs

XPath and XQuery:

- Navigation in XML structure

Relational Databases:

- Rules for relational table manipulation

Missing:

- Sound specification of IR tasks
- Rules for score propagation and correlation
- **Connection between the two**

Our Approach

* Three level DBMS:

■ Conceptual level:

- XPath+IR (NEXI)

■ Logical level:

- extended region algebra

■ Physical level:

- relational model

Intermediate level

- # Algebraic approach
 - XML navigation is supported
 - Ranking is a part of the algebra
- # Opportunities
 - Query rewriting and optimization
 - ... also for IR-like queries

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XML example

```
<article lang="en" date="10/02/04">  
  <title>Region algebra</title>  
  <bdy>  
    <sec>  
      <p>Structured documents ...</p>  
      <p> Text search ...</p>  
    </sec>  
    ...  
  </bdy>  
  ...  
</article>
```


XML example - Index (step 1)

```

<article0 lang1=2"en"3 date4=5"10/02/04"6>
  <title7>8Region9 algebra10</title>11
  <bdy>12
    <sec>13
      <p>14Structured15 documents16 ...</p>17
      <p>18Text19 search20 ...</p>21
    </sec>22
    ...
  </bdy>23
  ...
</article>24
  
```

- Start tag
- End tag
- Attribute name
- Attribute value
- Term

Indexed XML example (step 2)

Start	End	Name	Type
0	10034	article	node
1	2	lang	attr_name
2	2	en	attr_value
3	4	date	attr_name
4	4	10/02/04	attr_value
5	8	title	node
6	7	-	text
6	6	region	term
7	7	algebra	term
9	9876	bdy	node
10	576	sec	node
11	54	p	node
12	53	-	text
12	12	structured	term
13	13	documents	Term
...

The Storage of XML

Node table N

start	end	name	type
0	10034	article	node
5	8	title	node
6	8	-	text
9	9876	bdy	node
10	576	sec	node
11	54	p	node
12	53	-	text
...

Word table W

start	name
6	region
7	algebra
12	structured
13	documents
...	...

Attribute table A

start	owner	name	type
1	0	lang	name
2	0	en	value
3	0	date	name
4	0	10/02/04	value
...

Fragmentations

Horizontal

- XML node type

Vertical

- name and type of XML elements

Path-based

Not unified

Queries & Relational Algebra

* Bottleneck

- Descendant/ancestor step
- Join and projection combination

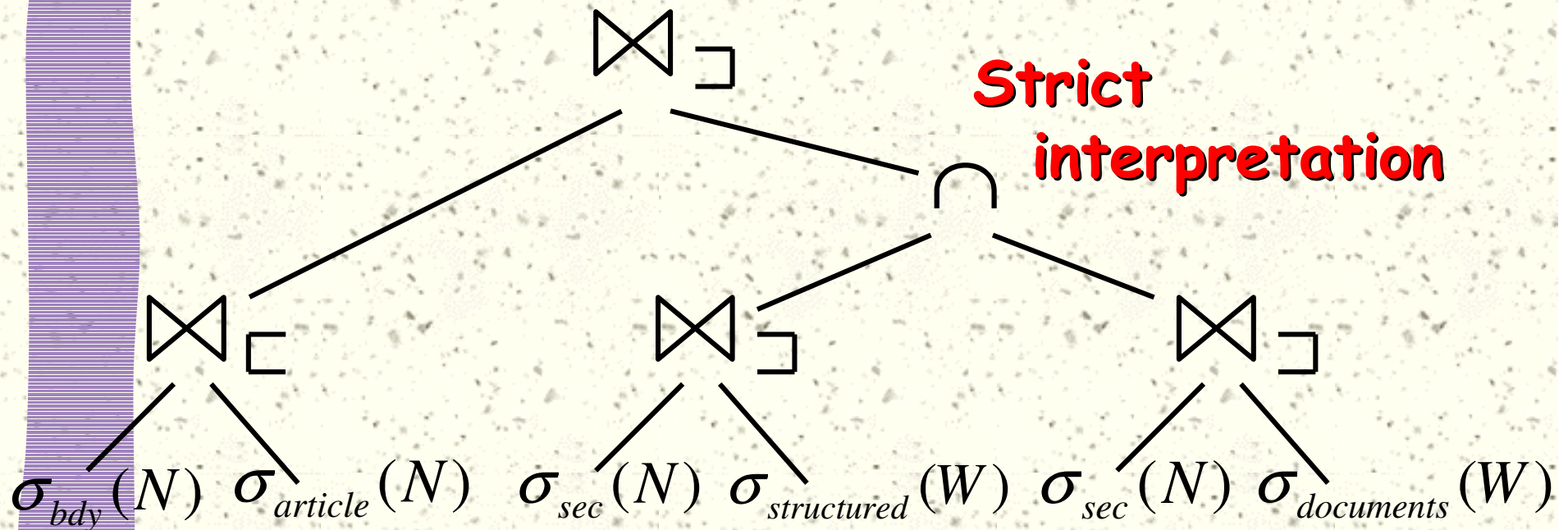
$$R = \pi_{start_2, end_2, name_2} (R_2 \bowtie_{start_2 > start_1, end_2 < end_1} R_1)$$

$$R = \pi_{start_2, end_2, name_2} (R_2 \bowtie_{start_2 < start_1, end_2 > end_1} R_1)$$

- "Containment join" (\bowtie_{\supseteq} and \bowtie_{\subseteq})

Example query

**//article//bdy[about(.//sec, structured)
and about(.//sec, documents)]**



Intermezzo: Logical Algebra

Relational algebra:

- New operators for IR-like queries
- Relational query plan highly dependant on relational storage
- Not XML (structure) aware

Logical Algebra

- Right level of abstraction for IR operators
- Data independence
- Query rewriting and optimization on logical level
- IR understanding and IR operator optimization

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Region Algebra

Domain $R = \{r \mid r = (s, e, n, t)\}$

Operators

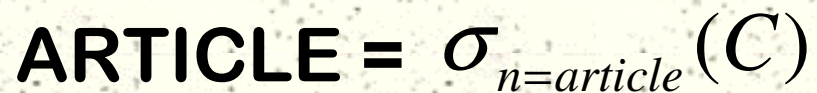
■ select $\sigma_{n=name}(R)$

■ containing \sqsupset

■ contained by \sqsubset

■ intersection \sqcap

■ union \sqcup



Logical query plan

**//article//bdy[about(., region) and about(., algebra)]
[about(./sec, XML) //p[about(., information) and
about(., retrieval)]**

R1 = BDY \sqcap ARTICLE

R2 = ((R1 \sqcap REGION) \sqcap (R1 \sqcap ALGEBRA)) \sqcap (SEC \sqcap XML)

R3 = P \sqcap R2

R4 = (R3 \sqcap INFORMATION) \sqcap (R3 \sqcap RETRIEVAL)

ARTICLE = $\sigma_{n=article}(C)$

Algebra Operator Properties (1)

* Regular

- Identity $\{(\sqcap, C), (\sqcup, \emptyset)\}$
- Commutativity $\{\sqcap, \sqcup\}$
- Associativity $\{\sqcap, \sqcup\}$
- Distributivity $\{(\sqsupset, \sqcup), (\sqsubset, \sqcup), (\sqcap, \sqcup), (\sqcup, \sqcap)\}$

Algebra Operator Properties (2)

Special cases

■ $op1 = \{\sqcup, \sqcap\}; op2 = \{\sqcup, \sqcap\}$

1

$$(R1 \ op1 \ R2) \ op2 \ R3 = (R1 \ op2 \ R3) \ op1 \ R2$$

$$(R1 \ op1 \ R2) \ op2 \ R3 = (R1 \ op1 \ R2) \sqcap (R1 \ op2 \ R3)$$

2

■ $op1 = \{\sqcap, \sqcup\}; op2 = \{\sqcup, \sqcap\}$

$$(R1 \ op1 \ R2) \ op2 \ R3 = (R1 \ op2 \ R3) \ op1 (R2 \ op2 \ R3)$$

3

Properties in action (1)

$(BDY \sqsubseteq ARTICLE) \sqsubseteq ((SEC \sqsubseteq STRUCTURED) \sqsubseteq (SEC \sqsubseteq DOCUMENTS))$

1

$(BDY \sqsubseteq ((SEC \sqsubseteq STRUCTURED) \sqsubseteq (SEC \sqsubseteq DOCUMENTS))) \sqsubseteq ARTICLE$

2

$(BDY \sqsubseteq ((SEC \sqsubseteq STRUCTURED) \sqsubseteq DOCUMENTS)) \sqsubseteq ARTICLE$

1

$(BDY \sqsubseteq ((SEC \sqsubseteq DOCUMENTS) \sqsubseteq STRUCTURED)) \sqsubseteq ARTICLE$

Properties in action (2)

$((\text{ARTICLE} \sqsubseteq \text{REGION}) \sqcap (\text{ARTICLE} \sqsubseteq \text{ALGEBRA}))$
 $\sqsubseteq (\text{SEC} \sqsubseteq \text{XML})$

3

$((\text{ARTICLE} \sqsubseteq \text{REGION}) \sqsubseteq (\text{SEC} \sqsubseteq \text{XML})) \sqcap$
 $((\text{ARTICLE} \sqsubseteq \text{ALGEBRA}) \sqsubseteq (\text{SEC} \sqsubseteq \text{XML}))$

1

$((\text{ARTICLE} \sqsubseteq (\text{SEC} \sqsubseteq \text{XML})) \sqsubseteq \text{REGION}) \sqcap$
 $((\text{ARTICLE} \sqsubseteq (\text{SEC} \sqsubseteq \text{XML})) \sqsubseteq \text{ALGEBRA})$

2

$((\text{ARTICLE} \sqsubseteq (\text{SEC} \sqsubseteq \text{XML})) \sqsubseteq \text{REGION}) \sqsubseteq \text{ALGEBRA}$

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Scoring in region algebra

Domain $R = \{r \mid r = (s, e, n, t, p)\}$

New operators

- ranked containing \supseteq_p
- ranked contained by \sqsubseteq_p
- ranked intersection \sqcap_p
- ranked union \sqcup_p

Scoring operators

$$\# R1 \sqsupset_p R2 \quad p = p_1 \bullet f_{\supset}(r_1, R_2)$$

$$\# R1 \sqsubset_p R2 \quad p = p_1 \bullet f_{\subset}(r_1, R_2)$$

$$\# R1 \sqcap_p R2 \quad p = p_1 \otimes p_2$$

$$\# R1 \sqcup_p R2 \quad p = p_1 \oplus p_2$$

Scoring functions and operators

Scoring functions:

- ▀ structural relation
- ▀ score values

$$f_{\supset}(r, R)$$

$$f_{\subset}(r, R)$$

Abstract operators:

- ▀ "and" expression
- ▀ "or" expression

$$\otimes = \{\bullet, \min, \dots\}$$

$$\oplus = \{+, \max, \dots\}$$

Complex scoring functions

$$f_{\subseteq}(r, R) = \sum_{\bar{r} \in R \subset R'} (g_{\subseteq}(\bar{r}, r) \bullet \bar{p})$$

$$f_{\supset}(r, R) = \sum_{\bar{r} \in R \subset R'} (g_{\supset}(\bar{r}, r) \bullet \bar{p})$$

$$R' = \{r\}$$

**Possible imple-
mentation of \mathbf{g}**

$$g_{\supset}(\bar{r}, r) = \frac{\text{size}(\bar{r})}{\text{size}(r)}$$

$$g_{\subseteq}(\bar{r}, r) = \frac{\text{size}(\bar{r})}{\sum_{\bar{r}} \text{size}(r)}$$

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Properties of scoring operators

$$(R \sqcap_p C) = (C \sqcap_p R)$$

$$p * 1 = 1 * p = p$$

$$(R \sqcup_p \emptyset) = (\emptyset \sqcup_p R)$$

$$p + 0 = 0 + p = p$$

$$\forall r \in R$$

$$R1 \sqcap_p (R2 \sqcup_p R3) = (R1 \sqcap_p R2) \sqcup_p (R1 \sqcap_p R3)$$

$$(R1 \text{ op1 } R2) \text{ op2 } R3 = (R1 \text{ op2 } R3) \text{ op1 } R2$$

$$\text{op1} = \{\sqcap_p, \sqcup_p\} \quad \text{op2} = \{\sqcap_p, \sqcup_p\}$$

$$p = (p_1 \bullet f(r_1, R_2)) \bullet f(r_1, R_3) = (p_1 \bullet f(r_1, R_3)) \bullet f(r_1, R_2)$$

only if $f(r, R) = f(s, n, t, R)$

... conditional properties(1) ...

$$(R1 \text{ op1 } R2) \text{ op2 } R3 = (R1 \text{ op1 } R2) \sqcap_p (R1 \text{ op2 } R3)$$

$$\text{op1} = \{\sqcap_p, \sqcup_p\} \quad \text{op2} = \{\sqcap_p, \sqcup_p\}$$

$$p = (1 \bullet \sum_{\bar{r}} (g(\bar{r}, r_1)) \bullet p_2) \bullet \sum_{\bar{r}} (g(\bar{r}, r_1)) \bullet p_3$$

$$= (1 \bullet \sum_{\bar{r}} (g(\bar{r}, r_1)) \bullet p_2)) \bullet (1 \bullet \sum_{\bar{r}} (g(\bar{r}, r_1)) \bullet p_3)$$

Conditional properties - example

$((P \sqsubseteq_p \text{SEC}) \sqsubseteq_p \text{INFORMATION}) \sqsubseteq_p$

$((P \sqsubseteq_p \text{SEC}) \sqsubseteq_p \text{RETRIEVAL})$



$((P \sqsubseteq_p \text{INFORMATION}) \sqsubseteq_p (P \sqsubseteq_p \text{RETRIEVAL}))$
 $\sqsubseteq_p \text{SEC}$

4

$((P \sqsubseteq_p \text{INFORMATION}) \sqsubseteq_p \text{RETRIEVAL})$
 $\sqsubseteq_p \text{SEC}$

Conditional properties (2)

$$\text{op1} = \{\sqcap_p, \sqcup_p\}$$

~~$$(R1 \sqcap_p R2) \text{ op1 } R3 = (R1 \text{ op1 } R2) \sqcap_p (R2 \text{ op1 } R3)$$~~

~~$$(p_1 \bullet p_2) \bullet f(r_{1,2}, R_3) = (p_1 \bullet f(r_{1,2}, R_3)) \bullet (p_2 \bullet f(r_{1,2}, R_3))$$~~

$$(R1 \sqcup_p R2) \text{ op1 } R3 = (R1 \text{ op1 } R2) \sqcup_p (R2 \text{ op1 } R3)$$

$$(p_1 + p_2) \bullet f(r_{1,2}, R_3) = (p_1 \bullet f(r_{1,2}, R_3)) + (p_2 \bullet f(r_{1,2}, R_3))$$

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Conclusions

- # Problem: Execution of IR-like queries over XML documents stored in relational database
- # Usefulness of intermediate logical level based on region algebra (with score computation)
 - Data independence between levels
 - Right level of abstraction (understanding IR)
 - Opportunities for query optimization on logical level (including ranking operators)

... still to come

- # Experimental evaluation: benefits of intermediate logical level
- # The definition of score operators
=> operator properties
- # Usage of different retrieval models
- # Theoretical foundation for the definition of score operators