

Arjun Kahlon

Assignment NP- Completeness

1.

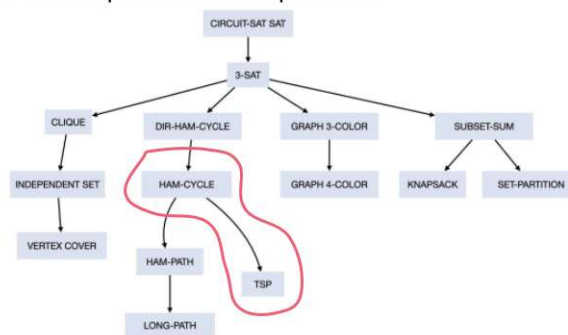
## Q1

First, let's show the problem is in NP

Given a graph  $G$ , paths of TSP and the number  $k$ , we can check for all the TSP paths against graph  $G$  to verify whether all the vertices are included and the cost is no larger than  $k$ . This check can be done in polynomial time.

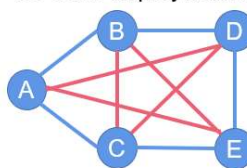
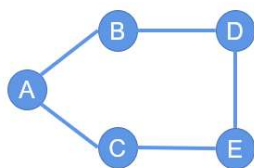
Second, let's show the problem is in NP-hard. Pick one known NP-hard problem and reduce it to our problem.

Let's pick the Hamiltonian cycle since that NP-hard problem is similar to our problem and it is easy to reduce from a similar problem to our problem.



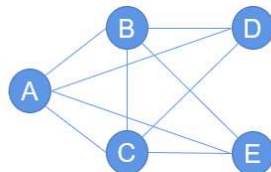
## Q1- cont'

Let's take a hamiltonian cycle.



For reduction, let's turn it to a graph  $G$  with all nodes connecting others. This operation needs to do  $(V-1+0)*V/2-V$  times so it can be done in polynomial time.

reference



A's round: 4 edges  
B's round: 3 edges  
C's round: 2 edges  
D's round: 1 edge  
E's round: 0 edge

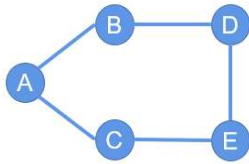
suppose we have  $V$  nodes

1st round:  $V-1$  edges  
2nd round:  $V-2$  edges  
...  
...  
last round: 0 edge

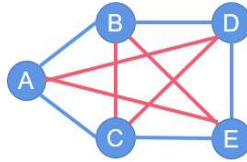
$$(v-1)+(v-2)+(v-3) + \dots + 0 = [(v-1)+0]*v/2$$

## Q1- cont'

Let's take a Hamiltonian cycle.



For reduction, let's turn it to a graph  $G'$  with all nodes connecting others. This operation needs to do  $(V-1+0)*V/2-V$  times so it can be done in polynomial time.



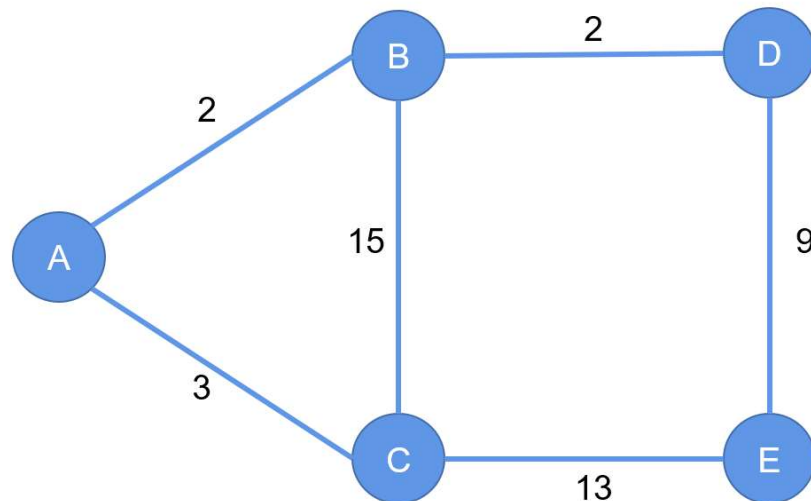
Suppose we have the 'magic polynomial time algorithm' for our TSP problem. From the graph  $G'$ , we can easily remove the edges that are not included in the cost of  $k$ . By doing that, we get a Hamiltonian cycle. However we know Hamiltonian cycle is NP-Hard, so we cannot have that 'magic polynomial time algorithm'. Therefore, our problem is NP-Hard.

Now we know the problem is in both NP and NP-Hard, so it is NP-Complete.

2.

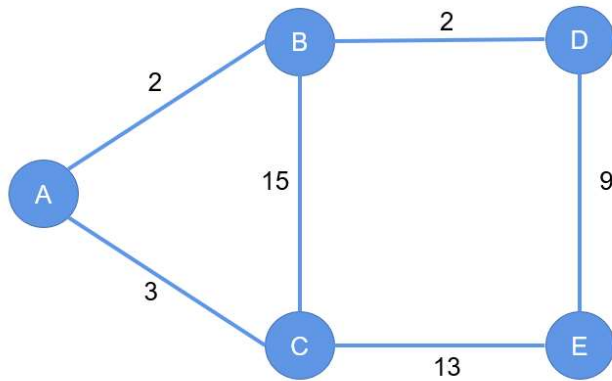
a)

## Q2 - a



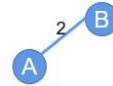
b)

## Q2 - b

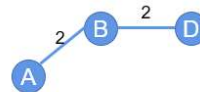


Let's use Prim's algorithm to get the MST first.

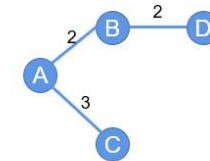
Step 1- Start from A and pick B.



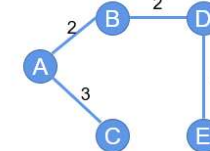
Step 2- Based on Step 1 pick D.



Step 3- Based on Step 2 pick C.

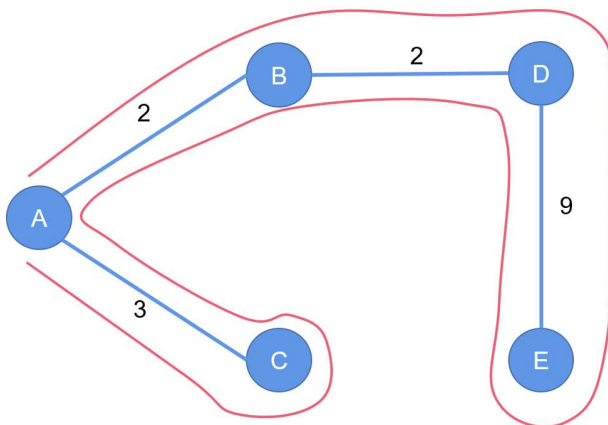


Step 4- Based on Step 3 pick E.



## Q2 - b cont'

MST



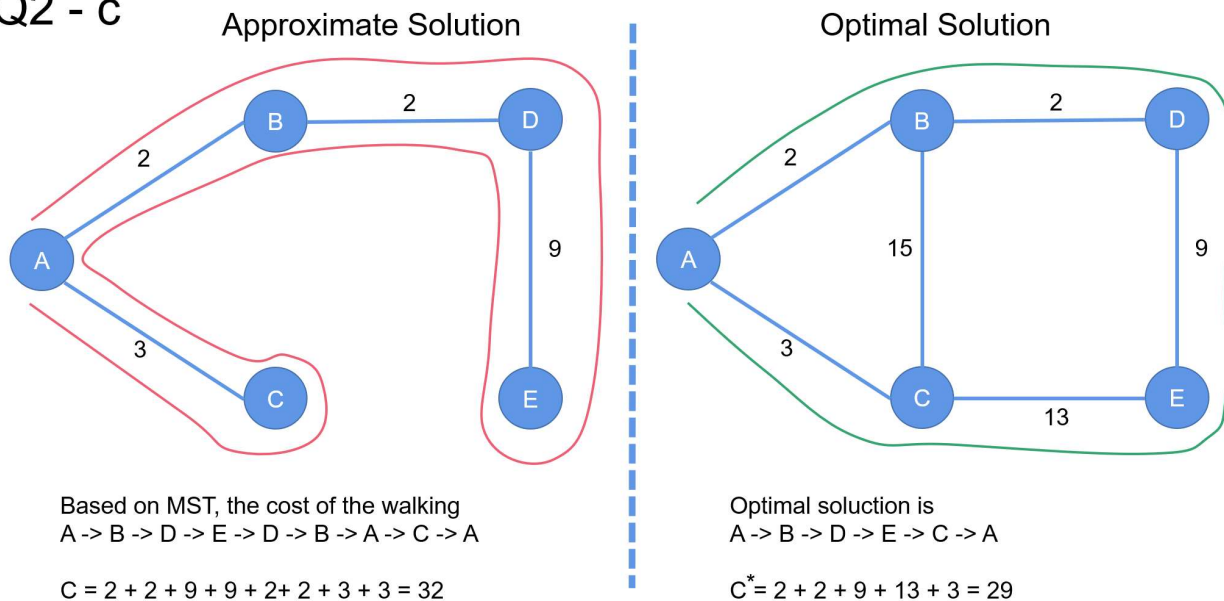
Let's walk based on the MST following below path.

A -> B -> D -> E -> D -> B -> A -> C -> A

So above is the approximate solution.

c)

## Q2 - c



As a result, we get accuracy ration  $\max(32/29, 29/32) = 1.1034 = 1.1$

d)

## Q2 - d

approx-nearest-neighbour-tsp(G):

  solution = {}

  picked\_nodes = {}

  remaining\_nodes = All nodes in G

  starting\_point\_node = randomly pick node from remaining\_nodes

  picked\_nodes.add starting\_point\_node

  remaining\_nodes.remove starting\_point\_node

  while (remaining\_nodes is not empty):

    nearest\_neighbour = pick the nearest node to picked\_nodes from remaining\_nodes

    solution.add path to nearest\_neighbour

    picked\_nodes.add nearest\_neighbour

    remaining\_nodes.remove nearest\_neighbour

  return solution

## Debriefing

1. I found this assignment to be quite difficult and spent 15-20 hours to fully complete it.
2. Overall, I would rate it as difficult as I found the material hard to conceptually grasp. I will need to spend more time reviewing this weekend.
3. I'd say I understand about 60% of the material. I hope to gain a better understanding by the end of this week.
4. I would appreciate a webinar on this material that would supplement the lessons from the Modules.