SSN COLLEGE OF ENGINEERING, KALAVAKKAM Department of Computer Science and Engineering UCS1503 - Theory of Computation

Tutorial

Due Date 08.10.2022

1. Consider the grammar $G = (\{S, A\}, \{a,b\}, P,S)$, where P consists of $S \rightarrow aAS \mid b$

S ZaAS | U

 $A \rightarrow SbA \mid ba$

Write the derivation and draw its equivalent parse tree for w = abbbab

- 2. If G is a grammar $S \rightarrow SbS \mid a$ prove that G is ambiguous.
- 3. Consider the grammar $S \rightarrow aS \mid aSbS \mid \epsilon$. This grammar is ambiguous. Show that the string aab has two
 - (a) Parse trees (b) Leftmost derivations (c) Rightmost derivations
- 4. For the grammar

 $S \rightarrow A1B$

 $A \rightarrow 0A \mid \epsilon$

 $B \rightarrow 0B|1B| \epsilon$

Give left most and right most derivation for the string 00101.

- 5. Construct CFG to generate $\{a^nb^n \mid n \in Z^+\}$
- 6. Consider the alphabet $\Sigma = \{a, b, (,), +, *, ., \epsilon \}$. Construct a context free grammar that generates all strings in Σ^* that are regular expressions over the alphabet $\{a,b\}$.
- 7. Convert the given CFG to CNF

 $S \rightarrow a \mid aA \mid B$

 $A \rightarrow aBB \mid \epsilon$

 $B \rightarrow Aa \mid b$

8. Convert the given CFG to GNF

 $S \rightarrow XB \mid AA$

 $A \rightarrow a \mid BA \mid AB$

 $B \to b$

 $X \rightarrow a$

- 9. Construct PDA
 - (a) L= $\{a^nb^mc^n/m,n>=1\}$ Acceptance by empting the stack
 - (b) L={ $wcw^R / w \in (a/b)^*$, $c \in a/b/epsilon$ } Acceptance by reaching the final state
 - (c) L= $\{a^nb^mc^md^n/m,n>=1\}$ Acceptance by empting the stack and reaching the final state