

# IE 555 – Programming for Analytics

## Simulated Annealing Report

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### Flow of the code

The flow of the code is represented in Figure 1: Flowchart of the travelling salesman problem. Initially we import all the libraries that would be required for our code to function. We then initialize the values of initial temperature, iteration, delta value, final temperature and cut off time. The inputs like locations folder, objective value, nearest neighbor, integer programming and simulated annealing selection type, type of map plot are captured from the command lines. Using the locations folder we import the values in the csv folder of Problem\_3.csv or Problem\_25.csv. We create a node using the class 'make\_node' and store all the details about the nodes. We use the function genTravelMatrix to get the distance and time data frame among the nodes from Mapquest using valid mapquest key. Now we solve the TSP problem and depending on our selection for the algorithm/s we have chosen to get the output. The selection goes to the respective chosen function and gives the tour and its cost.

The flow in Simulate Annealing function is described in Figure 2: Flowchart of the simulated annealing. In this function we receive all the required input like the data frames and the initialized variables required for simulated annealing. Initialize the start time and all the lists which would eventually store the points which has bad values and good values to plot. Call the nearest neighbor function to get its solution. Then we start with the Phase I of simulated annealing algorithm where we set the current tour and best tour as the nearest neighbor tour. Phase II starts with the iteration where we append in a list the Z values. 5 subtours reversals are generated using a while loop. We generate the subtour reversal by randomly selecting the starting and the end point and then reversing the selected part to get the tour. Depending on the objective type the cost is calculated and checked if it's the minimum cost among the 5. Then we check if the subtour cost is less than the current cost, if yes, we change the current tour to the reversed subtour, if no, then we check that if the reversed subtour is acceptable. The good solution, the bad solution accepted and the bad solution rejected is appended in lists for further plotting. We check if the current solution is better than best solution, then we change the best solution with the current. We finally check the breaking condition that if we have reached the final temperature or reached the cut-off time. We finally plot all the points and generate a plot to compare the cost value along with time.

After finding all the solutions for the respective algorithms we use Folium to plot the tours. We first check the depots and mark it in red. According to the given input we select if we need to plot the map turn by turn lines or with polylines. We use the cost values in each optimization technique to determine its presence and hence plot the tours. Finally we print the tours and their respective costs as output.

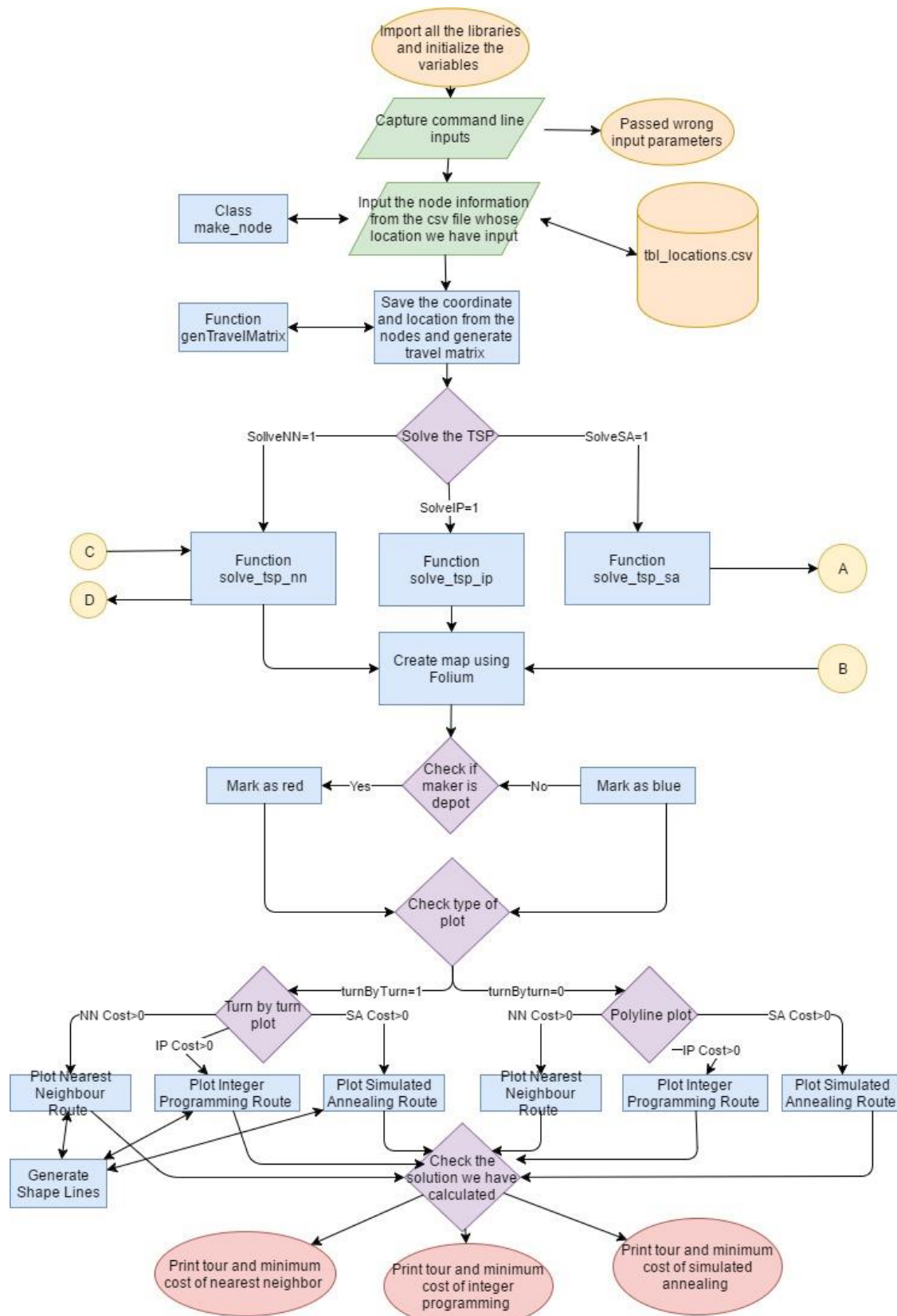


Figure 1: Flowchart of the Traveling Salesman Problem

Figure 2: Flowchart of the Simulated Annealing

The cost function for the distance table is as follows

**For Initial Temperature=200**

Delta\Iteration	500	1000	5000	10000	50000
0.5	70.503	65.929	67.456	69.392	66.443
0.1	72.939	66.655	68.613	37.825	65.938
0.05	71.134	69.363	66.824	68.36	69.647
0.01	68.41	70.33	70.33	68.594	60.473
0.005	69.393	67.456	71.061	62.942	61.252
0.001	67.84	70.538	65.966	67.253	65.614

Table 1: Initial Temperature 200

**For Initial Temperature=500**

Delta\Iteration	500	1000	5000	10000	50000
0.5	70.33	69.509	69.509	74.591	70.443
0.1	69.603	71.576	71.576	67.115	67.413
0.05	72.576	70.33	70.33	65.693	67.29
0.01	69.363	70.94	70.94	69.809	64.406
0.005	64.298	67.98	67.98	70.069	65.845
0.001	71.852	70.475	70.475	67.571	64.629

Table 2: Initial Temperature 500

**For Initial Temperature=1000**

Delta\Iteration	500	1000	5000	10000	50000
0.5	70.055	67.571	72.945	69.029	71.966
0.1	71.218	68.502	67.736	64.983	69.921
0.05	71.64	69.393	73.74	69.208	69.481
0.01	74.591	73.861	68.964	67.838	64.843
0.005	69.754	69.454	70.333	66.461	65.791
0.001	67.429	73.773	71.143	69.032	67.633

Table 3: Initial Temperature 1000

From the above observations we found the optimum value of 60.473 at Initial Temperature=200, Final temperature=0, delta=0.01 and iteration=50,000.

The cost function for the time table is as follows

**For Initial Temperature=200**

Delta\Iteration	500	1000	5000	10000	50000
0.5	8327	8327	8327	8248	8327
0.1	8327	8327	8165	8314	8327
0.05	8136	8248	8178	8322	8327
0.01	8327	8327	8050	8143	8054
0.005	8327	8327	8327	8089	8054
0.001	8274	8327	8327	8054	8140

Table 4: Initial Temperature 200

#### For Initial Temperature=500

Delta\Iteration	500	1000	5000	10000	50000
0.5	8327	8327	8327	8327	8323
0.1	8327	8327	8327	8193	8137
0.05	8327	8327	8327	8327	8179
0.01	8327	8327	8327	8327	8047
0.005	8327	8327	8327	8327	8226
0.001	8327	8327	8327	8327	8327

Table 5: Initial Temperature 500

#### For Initial Temperature=1000

Delta\Iteration	500	1000	5000	10000	50000
0.5	8327	8327	8327	8327	8327
0.1	8327	8327	8327	8327	8321
0.05	8327	8327	8327	8327	8080
0.01	8327	8327	8327	8327	8327
0.005	8327	8327	8327	8327	8327
0.001	8327	8327	8327	8327	8327

Table 6: Initial Temperature 1000

**From the above observations we found the optimum value of 8054 at Initial Temperature=200, Final temperature=0, delta=0.001/0.005 and iteration=10,000/50,000 respectively.**

From the above observations we found that more the iterations, we have greater chances of finding a good value. For delta we can say that it has an optimum value around the value 0.05 and 0.01. We see that in lesser values and greater values of delta, the cost value is increasing. Delta also determines the number of iterations. We can also observe in Table 1 that T=200 is giving more better solution than the other two in case of distance. By seeing the values in the temperature we can say that there is a possibility to get better distance values if we decrease the temperature. In case of objective value of temperature we can see from Table 4, Table 5 and Table 6, that higher temperature and lower iterations are not giving a better solution. Some particular combinations, which would work poorly, will be having low temperature and high delta. This would not allow much iterations to find the good solutions as temperature will decrease very soon.

#### Running the heuristic for 10 times

Sl No	Iteration	Objective Value	Runtime(seconds)
1	50000	67.099	35.1769998074
2	50000	63.609	35.2909998894
3	50000	67.12	35.2550001144
4	50000	59.431	34.9769999981
5	50000	59.81	35.2829999924
6	50000	64.746	35.4390001297
7	50000	66.89	34.8800001144
8	50000	62.728	34.9659998417

9	50000	65.623	34.9679999352
10	50000	59.71	35.2149999142

Table 4: Observation table for 10 runs

We don't get the same solution each time because each time we run the program the subtour reversals are generated randomly. The random subtour reversals have different cost values at different time of the iteration causing selection of different current solution, which would ultimately lead to the best solution. It also randomly accepts values whose probability decrease with the decrease in the best solution. In the above table we can see that the best solution ranges from 59.431 to 67.12. The runtime also changes which depends on the time taken to iterate through the algorithm. Thus we can conclude that as every time the code is run we randomly generate subtours and randomly accept current solution which are not good, hence we randomly generate outputs in each run.

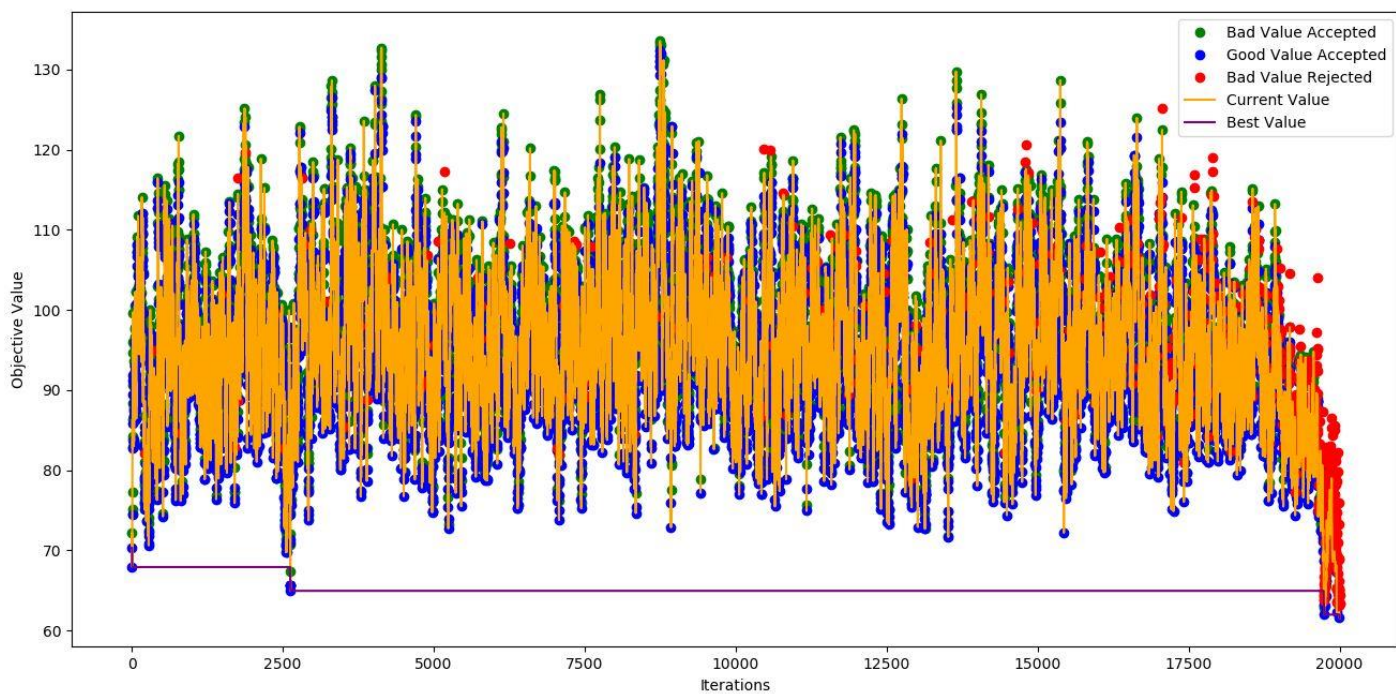


Figure 3: Simulated Annealing heuristic over time

We can refer to Figure 3: Simulated Annealing heuristic over time and infer some interesting observations.

Figure 3 depicts the flow of the objective value of the heuristics along with the iterations. In the figure there are two lines: the line in orange is the current value while the line in purple is the best value. There are 3 dot markers to represent 3 objective values: The blue dot shows the good values which were accepted, the green dot shows the bad values which were accepted and the red dot shows the bad values which were rejected. Thus in the figure we can observe how with increase in iterations the we are getting the best value in the last. In the figure we can also see that as the iteration is increasing we are rejecting bad values more than we are accepting the bad ones.

As the green dots indicate the bad value accepted we see that early in the process there are many bad values that are accepted as compared to the bad values that are rejected (red dots). As the temperature cools we can observe that the red dots increase significantly and hardly we can see any green dots when the temperature reaches the final temperature. We can also observe the heuristic is converging to a locally optimal solution as the current value is quite less near the final temperature.

### Comparison between the solution of Gurobi and Mixed Integer Linear Programming

We executed the code using the command “python solve\_tsp\_3.py practice\_25 1 0 1 1 120 1” in the console to run the MILP and SA heuristics.

```
Runtime of SA: 34 seconds
```

Figure 4: SA Runtime

```
Explored 1449 nodes (13376 simplex iterations) in 1.51 seconds
```

Figure 5: MILP(Gurobi) Runtime

Comparing Figure:4 and Figure:5 we can observe that using gurobi the runtime is a way more less than the runtime used in Simulate Annealing. The runtime of Gurobi is about 4.41% of the time taken by Simulated Annealing to reach its optimal value.

```
IP Route:
[0, 24, 5, 10, 12, 1, 2, 8, 11, 4, 18, 22, 16, 3, 17, 7, 19, 23, 13, 6, 15, 14, 21, 9, 20, 0]
IP 'cost':
56.444

SA Route:
[0, 24, 5, 10, 12, 1, 2, 20, 9, 7, 17, 3, 16, 22, 18, 4, 8, 11, 13, 23, 19, 14, 15, 6, 21, 0]
SA 'cost':
66.076
```

Figure 6: Route and objective values for distance

Figure 6 shows that the objective function value for MILP is much less than the cost of SA.

The runtime and the objective values clearly shows that Mixed Integer Linear Programming is a better choice to find the optimal solution than Simulated Annealing. The cost for MILP is about 15% less than the cost we get by Simulated Annealing.



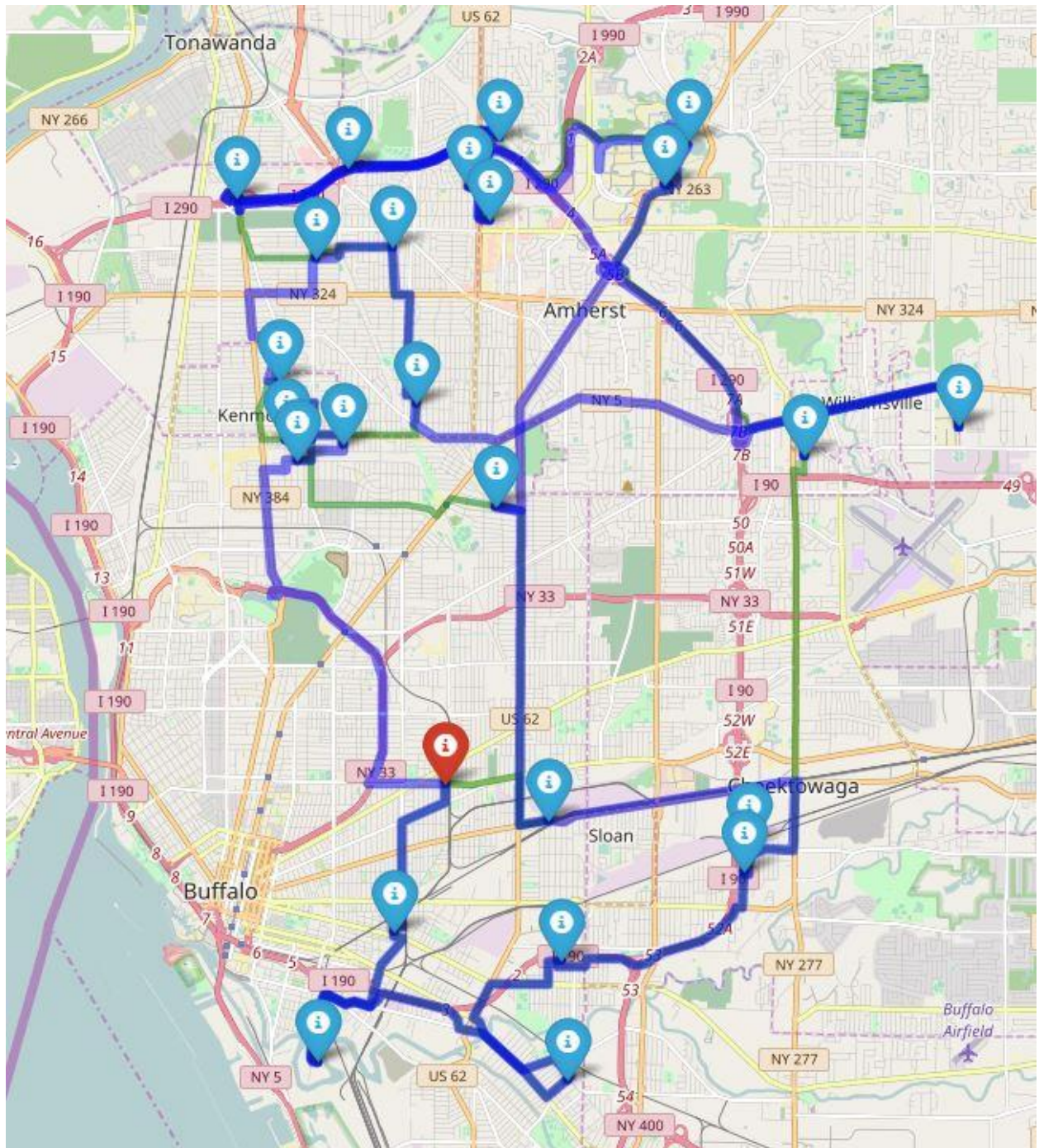


Figure 7: Map showing the route of SA in blue and MILP in green.



## Instructions for running the code

- The file name of the code is ARKACHOW\_code.py
- Make sure that the folder in which the file is present also has a folder named Problems which further has a folder named practice\_25 which contains a file named tbl\_locations.csv, it contains the location information.
- Open your terminal and go to the folder location where the file is present.
- Make sure you have the libraries random, matplotlib, numpy, time and math as additional libraries to run this code.
- Run the code along with the command line inputs.
- The input to be given in the command line while executing the code is given below
  - Locations Folder: practice\_25 (practice 3 will not work to generate 5 unique subtours)
  - Objective Type: 1- to minimize distance, 2- to minimize time
  - 1- to solve using nearest neighbor, 0- to not solve using nearest neighbor
  - 1- to solve using mixed integer linear programming, 0- to not solve using nearest neighbor
  - 1- to solve using simulated annealing, 0- to not solve using simulated annealing
  - Cut off time: 1- for no cut off time, 'value' in seconds for Gurobi cut off time in IP.
  - Map Route Display Type: 1- For detailed route, 0- Draw straight lines joining the points
- Sample line to be written in the terminal to run only simulated annealing- python solve\_tsp\_3.py practice\_25 1 0 0 1 120 1
- Code to run to compare MILP and SA- python solve\_tsp\_3.py practice\_25 1 0 1 1 120 1
- Code to run to get the optimum time from all the heuristics- python solve\_tsp\_3.py practice\_25 2 1 1 1 120 1