

## Game Theory -

Sandeep Kumar Gour

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Game is defined as an activity between two or more persons involving activities by each person according to a set of rules, at the end of which each person receives some benefit or satisfaction or suffers loss (negative benefit)

The set of rules defines the game. Going through the set of rules once by the participants defines a play.

## Characteristics of Game Theory-

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There can be various types of games. They can be classified on the basis of the following characteristics -

### (i) Chance of strategy -

If in a game, activities are determined by skill, it is said to be a game of strategy, if they are determined by chance, it is a game of chance.

In general, a game may involve game of strategy as well as a game of chance.

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In general, a game may involve game of strategy as well as a game of chance.

(ii) Number of persons -

A game is called an n-person game if the number of persons playing is n. The person means an individual or a group aiming at a particular objective.

(iii) Number of activities -

These may be finite or infinite.

(iv) Number of alternatives (choices) available to each person in a particular activity may also be finite or infinite. A finite game has a finite number of activities, each involving a finite number of alternatives, otherwise the game is said to be infinite.

(v) Information to the players about the past activities of other players is completely available, partly available or not available at all.

(vi) Payoff -

A quantitative measure of satisfaction a person gets

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(vi) Payoff -

A quantitative measure of satisfaction a person gets at the end of each play is called a payoff. It is a real-valued function of variables in the game.

## Basic Definitions -

### I. Competitive Game -

A competitive situation is called a competitive game if it has the following four properties -

- (i) There are finite number ( $n$ ) of competitors (called players).
- (ii) Each player has a list of finite number of possible activities.
- (iii) A play is said to occur when each player chooses one of his activities.
- (iv) Every combination of activities determines an outcome.

## 2. Zero-sum and Non-zero sum game - Sandeep Kumar Gour

Competitive games are classified according to the number of player involved, i.e as a two-person game, three-person game etc.

If the players make payments only to each other, i.e the loss of one is the gain of others and nothing comes from outside, the competitive game is said to be zero-sum.

A game which is not zero sum is called a Non-zero sum game.

B

Player A +10

number of player involved, i.e as a two-person, one person game etc.

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B  
player A +10

n players, payoff

$v_1 \ v_2 \ \dots \ v_n$

$p_1$

$p_2$

:

$p_n$

$$\boxed{\sum_{i=1}^n v_i = 0}$$

### 3. Strategy -

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A strategy for a given player is a set of rules that specifies which of the available course of action he should make at each play. This strategy may be of two kinds

#### (i) Pure strategy -

If a player knows exactly what the other player is going to do, a deterministic situation is obtained and objective function is to maximize the gain. Therefore, the pure strategy is a decision rule always to select a particular course of action.

#### (ii) Mixed strategy -

If a player is guessing as to which activity is to be selected by the other on any particular occasion, a probabilistic

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(ii) Mixed strategy -

If a player is guessing as to which activity is to be selected by the other on any particular occasion, a probabilistic situation is obtained and objective function is to maximise the expected gain.

Thus, the mixed strategy is a selection among pure strategies with fixed probabilities.

#### 4. Two person, zero sum (Rectangular) Games -

A game with only two players (say, player A and player B) is called a 'two person zero sum game' if the losses of one player are equivalent to the gains of the other, so that the sum of their net gains is zero.

Two person, zero sum games are also called rectangular games as these are usually represented by a payoff matrix in rectangular form.



## 5. Payoff matrix -

Suppose the player A has  $m$  activities and the player B has  $n$  activities. Then a payoff matrix can be formed by adopting the following rules.

- (i) Row designations for each matrix are activities available to player A.
- (ii) Column designations for each matrix are activities available to player B.
- (iii) Cell entry  $v_{ij}$  is the payment to player A in A's payoff matrix when A chooses the activity  $i$  and B chooses the activity  $j$ .
- (iv) In a 'sum, two person game', the cell entry in the player B's payoff matrix will be negative of the corresponding cell entry in the player A's payoff matrix.

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- (i) Row designations for each matrix are activities available to player A.
- (ii) Column designations for each matrix are activities available to player B.
- (iii) Cell entry ' $V_{ij}$ ' is the payment to player A in A's payoff matrix when A chooses the activity  $i$  and B chooses the activity  $j$ .
- (iv) With a 'zero sum, two person game', the cell entry in the player B's payoff matrix will be negative of the corresponding cell entry in the player A's payoff matrix.

player B has  $n$  activities. Then a payoff matrix can be formed by adopting the following rules.

- (i) Row designations for each matrix are activities available to player A.
- (ii) Column designations for each matrix are activities available to player B.
- (iii) Cell entry ' $v_{ij}$ ' is the payment to player A in A's payoff matrix when A chooses the activity  $i$  and B chooses the activity  $j$ .
- (iv) With a 'zero sum, two person game', the cell entry in the player B's payoff matrix will be negative of the corresponding cell entry ' $v_{ij}$ ' in the player A's payoff matrix so that sum of payoff matrices for player A and player B is ultimately zero.

player A's payoff matrix

		player B				
		1	2	3	...	n
player A	1	$v_{11}$	$v_{12}$	$\dots$	$v_{1n}$	
	2	$v_{21}$	$v_{22}$	$\dots$	$v_{2n}$	
	3	:				
	:	:				
	m	$v_{m1}$	$v_{m2}$	$\dots$	$v_{mn}$	

player B's payoff mat

		player B			
		1	2	...	n
player A	1	$-v_{11}$	$-v_{12}$	$\dots$	$-v_{1n}$
	2	$-v_{21}$	$-v_{22}$	$\dots$	$-v_{2n}$
	3	:	:		
	:	:			
	m	$-v_{m1}$	$-v_{m2}$	$\dots$	$-v_{mn}$

	1	2	3	...	$n$
1	$v_{11}$	$v_{12}$	$\dots$		$v_{1n}$
2	$v_{21}$	$v_{22}$	$\dots$		$v_{2n}$
3					
:					
$m$	$v_{m1}$	$v_{m2}$	$\dots$		$v_{mn}$

Eg: 2- coin

II, II

		Player B			
		H	T		
Player A	H	+1	-1	= 0	
	T	-1	+1		

		Player B			
		1	2	...	n
Player A	1	-v <sub>11</sub>	-v <sub>12</sub>	...	-v <sub>1n</sub>
	2	-v <sub>21</sub>	-v <sub>22</sub>	...	-v <sub>2n</sub>
	:	:			
	m	-v <sub>m1</sub>	-v <sub>m2</sub>	...	-v <sub>mn</sub>

### Maximin Criteria -

The player who is maximising his outcome or payoff finds out his minimum gains from each strategy (course of action) and selects the maximum value out of these minimum gains.

### Minimax Criteria -

In this criterion the minimising player determines the maximum loss from each strategy and then selects with minimum loss out of the maximum loss list.

### Saddle point -

A saddle point of a payoff matrix is the position of such an element in the payoff matrix which is minimum in its row and maximum in its column.

Mathematically :  $\exists$

## Minimax Criteria -

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In this criterion the minimising player determines the maximum loss from each strategy and then selects with minimum loss out of the maximum loss list.

## Saddle point -

A saddle point of a payoff matrix is the position of such an element in the payoff matrix which is minimum in its row and maximum in its column.

Mathematically, if a payoff matrix  $\{v_{ij}\}$  is such that

$$\max_i \left[ \min_j \{v_{ij}\} \right] = \min_j \left[ \max_i \{v_{ij}\} \right] = v_{rs} \text{ (say)}$$

then the matrix is said to have a saddle point  $(r, s)$

## Saddle point -

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A saddle point of a payoff matrix is the position of such an element in the payoff matrix which is minimum in its row and maximum in its column.

## Rules for determining a saddle point -

1. Select the minimum element of each row of the payoff matrix and mark them by 'O'
2. Select the greatest element of each column of the payoff matrix and mark them by 'I'
3. If there appears an element in the payoff matrix marked by 'O' and 'I' both, the position of that element is a saddle point of the payoff matrix.

matrix and mark them by  $\square$

3. If there appears an element in the payoff matrix marked by '0' and ' $\square$ ' both, the position of that element is a saddle point of the payoff matrix.

### Optimal strategy -

The course of action which maximizes the profit of a player or minimizes his loss is called an optimal strategy.

If the payoff matrix  $\{v_{ij}\}$  has the saddle point  $(r,s)$  then the players (A and B) are said to have  $r^{\text{th}}$  and  $s^{\text{th}}$  optimal strategies, respectively.

## Value of Game -

The payoff ( $v_{rs}$ ) at the saddle point ( $r, s$ ) is called the value of game and it is obviously equal to the maximin ( $\underline{v}$ ) and minimax value ( $\bar{v}$ ) of the game.

A game is said to be a fair game if  $\bar{v} = \underline{v} = 0$ .  
A game is said to be strictly determinable if  $\bar{v} = v = \underline{v}$



to the player A. find the optimal strategy, if any

		player B		
		I	II	III
player A		I	-3	6
		II	2	0

A → maximising  
B → minimising

Sol:-

		player B		
		I	II	III
player A		I	-3	6
		II	2	0
		III	5	-4

Row minimum

-3

0 maximin value (V)

-4

column maximum

5 0 6

Eg ②

Player B			
			Row minimum
			-4 ✓ (V)
Player A	I	3	-4
	II	-8	5
	III	6	-7

column  
maximum

6 ✓ 5 8  
minimax value  
(V)

Such game are said to be the games  
without saddle point.

## i) Analytical method -

A  $2 \times 2$  payoff matrix where there is no saddle point can be solved by Analytical method.

Given the matrix

$$\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}$$

Value of the game is

$$V = \frac{a_{11}a_{22} - a_{21}a_{12}}{(a_{11} + a_{22}) - (a_{12} + a_{21})}$$

with the coordinates

$$x_1 = \frac{a_{22} - a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})}, \quad x_2 = \frac{a_{11} - a_{12}}{(a_{11} + a_{22}) - (a_{12} + a_{21})}$$

$$\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}$$

value of the game is

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$$y_1 = \frac{a_{22} - a_{12}}{(a_{11} + a_{22}) - (a_{12} + a_{21})}, \quad y_2 = \frac{a_{11} - a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})}$$

Q: solve the payoff matrix -

$$\begin{array}{c} \text{player B} \\ \text{player A} \begin{bmatrix} 5 & 1 \\ 3 & 4 \end{bmatrix} \end{array}$$

Sol<sup>n</sup>:-

Sol :- value of the game

$$V = \frac{a_{11}a_{22} - a_{12}a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{5 \times 4 - 3 \times 1}{(5+4) - (1+3)} = \frac{17}{5}$$

with coordinates

$$x_1 = \frac{a_{22} - a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{4 - 3}{5} = \frac{1}{5}$$

$$x_2 = \frac{a_{11} - a_{12}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{5 - 1}{5} = \frac{4}{5}$$

$$y_1 = \frac{a_{22} - a_{12}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{4 - 1}{5} = \frac{3}{5}$$

$$y_2 = \frac{a_{11} - a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{5 - 3}{5} = \frac{2}{5}$$

$$\therefore \text{value of } V = \frac{a_{11}a_{22} - a_{12}a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{(5+4) - (1+3)}{(5+4) - (1+3)}$$

with coordinates

$$x_1 = \frac{a_{22} - a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{4-3}{5} = \frac{1}{5}$$

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$$y_1 = \frac{a_{22} - a_{12}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{4-1}{5} = \frac{3}{5}$$

$$y_2 = \frac{a_{11} - a_{21}}{(a_{11} + a_{22}) - (a_{12} + a_{21})} = \frac{5-3}{5} = \frac{2}{5}$$

optimal strategy of player A =  $(\frac{1}{5}, \frac{4}{5})$   
 B =  $(\frac{3}{5}, \frac{2}{5})$

~~then~~ Alternate procedure

Qui:

		Player B	
		1	1/5
Player A		4	4/5
		3	2
		3/5	2/5

value of game

$$\begin{aligned} v &= 5 \times \frac{1}{5} + 3 \times \frac{4}{5} \\ &= 1 + \frac{12}{5} \\ &= \frac{17}{5} \end{aligned} \quad \left| \begin{array}{l} \text{or} \\ = 5 \times \frac{3}{5} + 1 \times \frac{2}{5} \\ = 3 + \frac{2}{5} \\ = \frac{17}{5} \end{array} \right.$$

$$\begin{array}{cc} & \begin{matrix} 3 & 2 \end{matrix} \\ \begin{matrix} 3 \\ 5 \end{matrix} & \begin{matrix} 2/5 \\ 2/5 \end{matrix} \end{array}$$

value of game

$$\begin{aligned} V &= 5 \times \frac{1}{5} + 3 \times \frac{4}{5} \\ &= 1 + \frac{12}{5} \\ &= \frac{17}{5} \end{aligned} \quad \left| \begin{array}{l} \text{or} \\ = 5 \times \frac{3}{5} + 1 \times \frac{2}{5} \\ = 3 + \frac{2}{5} \\ = \frac{17}{5} \end{array} \right.$$

optimal strategy of player A =  $(\frac{1}{5}, \frac{4}{5})$

B =  $(\frac{3}{5}, \frac{2}{5})$

## Graphical method -

The graphical method is used to solve the games whose payoff matrix has

- (1) Two rows and n columns ( $2 \times n$ )
- (2) m rows and 2 columns ( $m \times 2$ )

### ① Algorithm for solving $2 \times n$ matrix games -

- ① Draw two vertical axes 1 unit apart. The two lines are  $x_1=0$ ,
- ② Take the points of the first row in the payoff matrix on the vertical line  $x_1=1$  and the points of the second row in the payoff matrix on the vertical lines  $x_1=0$
- ③ The point  $a_{ij}$  on axis  $x_1=1$  is then joined to the point  $a_{2j}$  on the axis  $x_1=0$  to give a straight line. Draw 'n' straight lines for  $j=1, 2, \dots, n$  and determine the highest point of the lower envelope obtained. This will be the value of the game.

(2) m rows and 2 columns ( $m \times 2$ )

### ① Algorithm for solving $2 \times n$ matrix games -

- ① Draw two vertical axes 1 unit apart. The two lines are  $x_1=0$
- ② Take the points of the first row in the payoff matrix on the vertical line  $x_1=1$  and the points of the second row in the payoff matrix on the vertical lines  $x_1=0$
- ③ The point  $a_{1j}$  on axis  $x_1=1$  is then joined to the point  $a_{2j}$  on the axis  $x_1=0$  to give a straight line. Draw 'n' straight lines for  $j=1, 2, \dots, n$  and determine the highest point of the lower envelope obtained. This will be the maximin point.
- ④ The two or more lines passing through the maximin point determines the required  $2 \times 2$  payoff matrix. This in turn gives the optimum solution by making use of analytical method.

Eg: solve by graphical method

player B

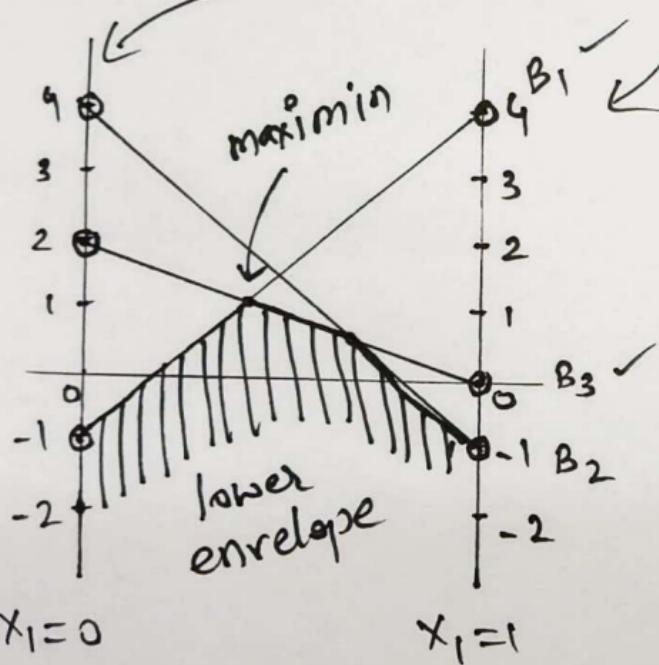
	$B_1$	$B_2$	$B_3$
$A_1$	4	-1	0
$A_2$	-1	4	2

player A

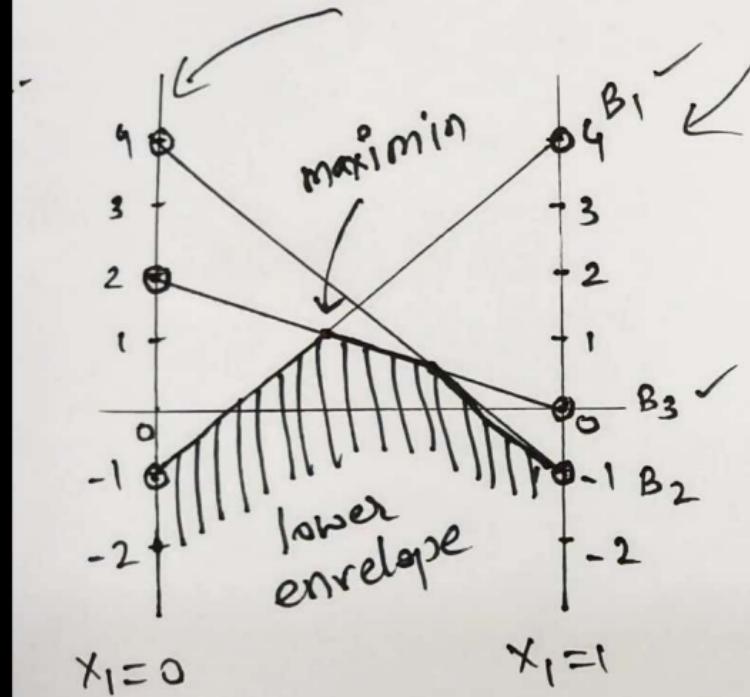
$\xrightarrow{2 \times 3}$

$\xrightarrow{2 \times 2}$

Sol:-



		player B			
		B <sub>1</sub>	B <sub>2</sub>	B <sub>3</sub>	
player A	A <sub>1</sub>	4	-1	0	$\frac{2+3}{2+3}$
	A <sub>2</sub>	-1	4	2	



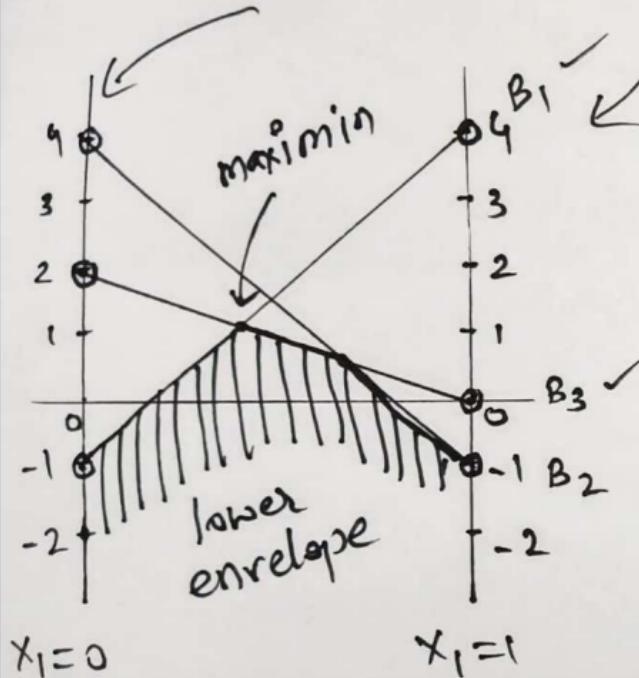
		B <sub>1</sub>	B <sub>3</sub>	
		4	0	$3 \frac{3}{7}$
A <sub>1</sub>	A <sub>2</sub>	-1	2	$4 \frac{4}{7}$
		2	5	

value of the game

$$\begin{aligned}
 v &= 4 \times \frac{3}{7} + (-1) \times \frac{4}{7} \\
 &= \frac{12}{7} - \frac{4}{7} \\
 &= \frac{8}{7}
 \end{aligned}$$

optimal strategy of player A =

$$\text{Player A} \quad \begin{array}{c} A_1 \\ A_2 \end{array} \left[ \begin{array}{ccc} 4 & -1 & 0 \\ -1 & 4 & 2 \end{array} \right] \rightarrow \begin{array}{c} A_2 \\ B_1 \\ B_2 \\ B_3 \end{array} \left[ \begin{array}{cc} -1 & 2 \\ 2 & 5 \\ \frac{2}{7} & \frac{8}{7} \end{array} \right] \leftarrow \frac{4}{7}$$



value of the game

$$\begin{aligned} v &= 4 \times \frac{3}{7} + (-1) \times \frac{4}{7} \\ &= \frac{12}{7} - \frac{4}{7} \\ &= \frac{8}{7} \end{aligned}$$

optimal strategy of player A =  $(\frac{3}{7}, \frac{4}{7})$   
 \_\_\_\_\_  
 B =  $(\frac{2}{7}, 0, \frac{5}{7})$

## (2) Algorithm for solving $m \times 2$ matrix games-

- (1) Draw two vertical axes 1 unit apart. The two lines are  $x_1=0$ ,  $x_1=1$
- (2) Take the points of the first column in the payoff matrix on the vertical line  $x_1=1$  and the points of the second column in the payoff matrix on the vertical line  $x_1=0$
- (3) The point  $a_{1j}$  on axis  $x_1=1$  is then joined to the point  $a_{2j}$  on the axis  $x_1=0$  to give a straight line. Draw 'n' straight lines for  $j=1, 2, \dots, n$  and determine the lowest point of the upper envelope obtained. This will be the minimax point.
- (4) The two or more lines passing through the minimax point determines the required  $2 \times 2$  payoff matrix.

(1) Draw two vertical axes 1 unit apart. The two lines are  
 $x_1=0$ ,  $x_1=1$

(2) Take the points of the first column in the payoff matrix on the vertical line  $x_1=1$  and the points of the second column in the payoff matrix on the vertical line  $x_1=0$

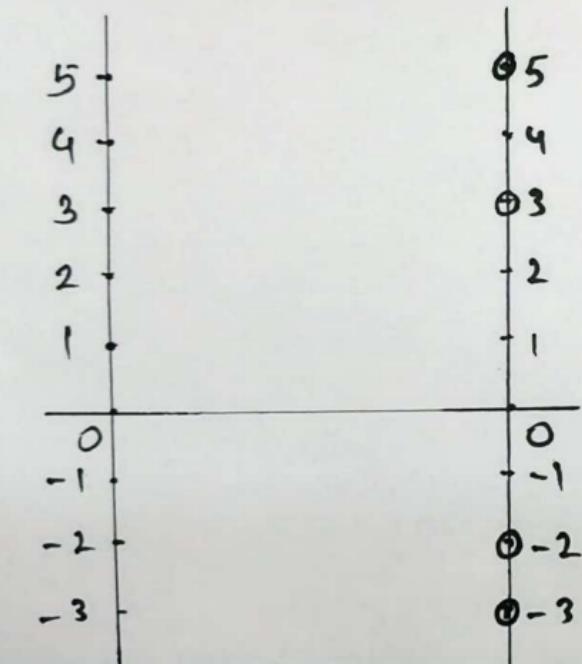
(3) The point  $a_{1j}$  on axis  $x_1=1$  is then joined to the point  $a_{2j}$  on the axis  $x_1=0$  to give a straight line. Draw 'n' straight lines for  $j=1, 2, \dots, n$  and determine the lowest point of the upper envelope obtained. This will be the minimax point.

(4) The two or more lines passing through the minimax point determines the required  $2 \times 2$  payoff matrix. This in turn gives the optimum solution by minimax.

$x_1=0$ ,  $x_1=1$

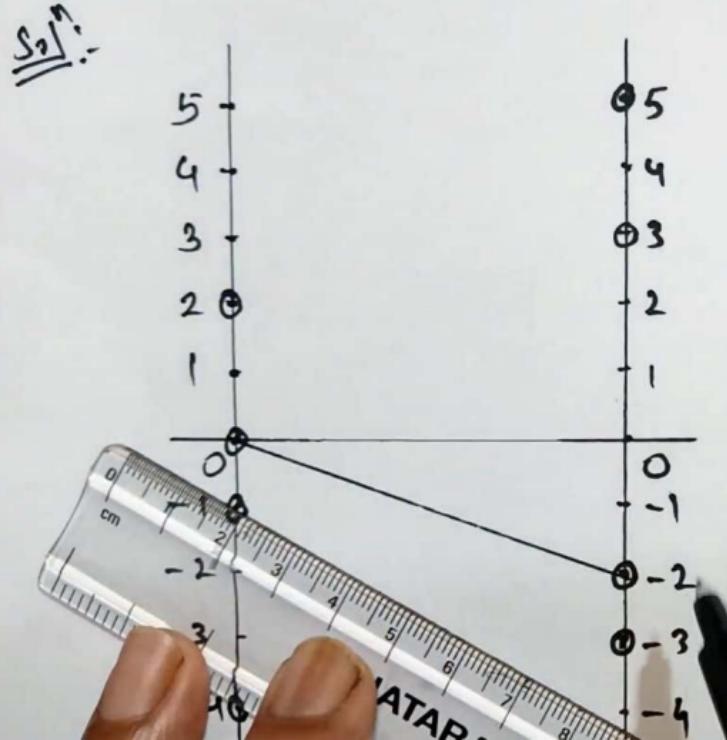
- (2) Take the points of the first column in the payoff matrix on the vertical line  $x_1=1$  and the points of the second column in the payoff matrix on the vertical line  $x_1=0$
- (3) The point  $a_{1j}$  on axis  $x_1=1$  is then joined to the point  $a_{2j}$  on the axis  $x_1=0$  to give a straight line. Draw 'n' straight lines for  $j=1, 2, \dots, n$  and determine the lowest point of the upper envelope obtained. This will be the minimax point.
- (4) The two or more lines passing through the minimax point determines the required  $2 \times 2$  payoff matrix. This in turn gives the optimum solution by making use of analytical method.

	$B_1$	$B_2$
$A_1$	-2	0
$A_2$	3	-1
$A_3$	-3	2
$A_4$	5	-4



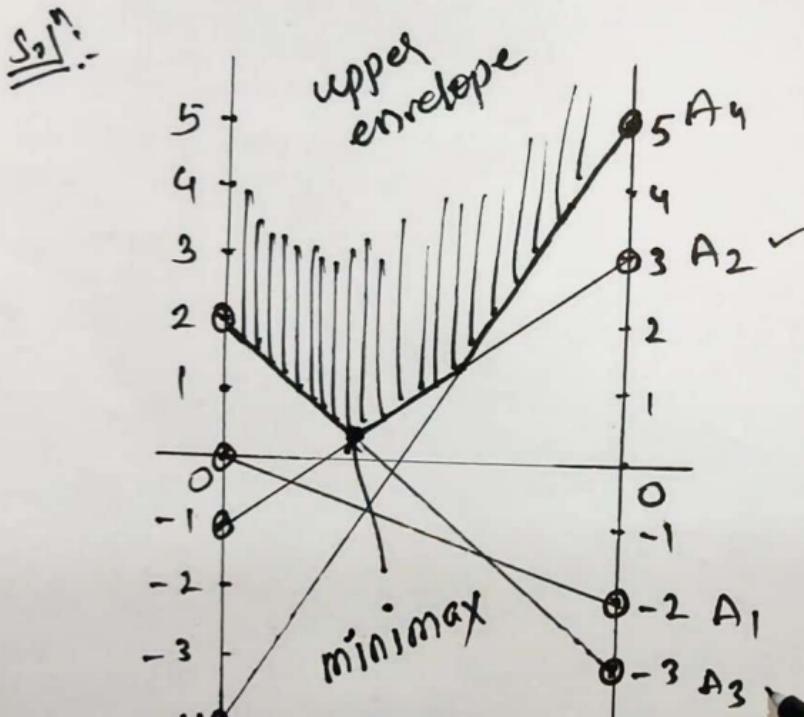
player A

$$\begin{array}{l}
 A_1 \\
 A_2 \\
 A_3 \\
 A_4
 \end{array}
 \left[ \begin{array}{cc}
 -2 & 0 \\
 3 & -1 \\
 -3 & 2 \\
 5 & -4
 \end{array} \right]$$



player A

	$B_1$	$B_2$
$A_1$	-2	0
$A_2$	3	-1
$A_3$	-3	2
$A_4$	5	-4



Eg:- solve by graphical method

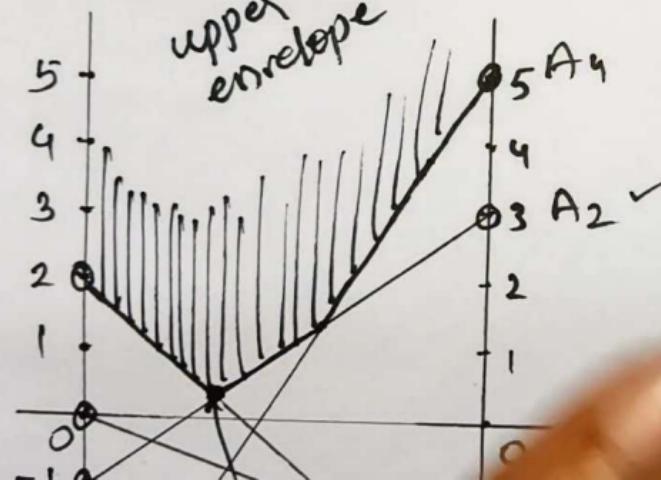
player B

~~m x 2~~

player A

$$\begin{array}{c} \text{player B} \\ \begin{array}{cc} B_1 & B_2 \end{array} \\ \begin{array}{c} A_1 \\ A_2 \\ A_3 \\ A_4 \end{array} \left[ \begin{array}{cc} -2 & 0 \\ 3 & -1 \\ -3 & 2 \\ 5 & -4 \end{array} \right] \end{array} \rightarrow \begin{array}{cc} B_1 & B_2 \\ A_2 & \left[ \begin{array}{cc} 3 & -1 \end{array} \right] \\ A_3 & \left[ \begin{array}{cc} -3 & 2 \end{array} \right] \end{array}$$

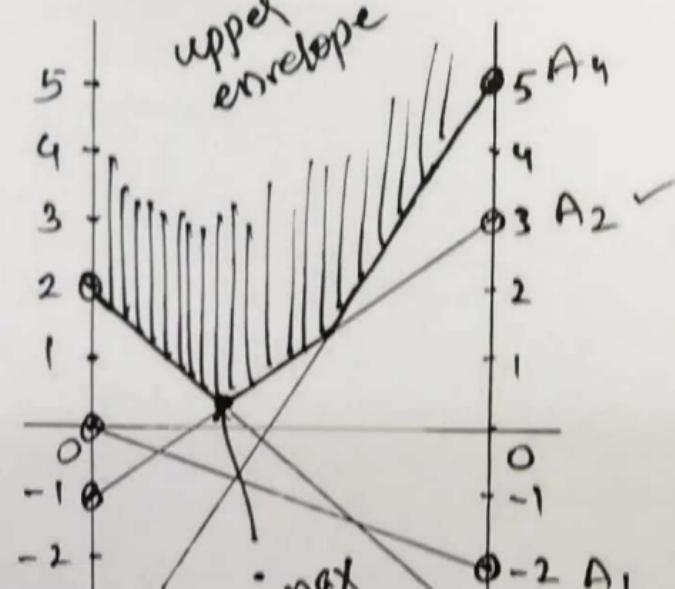
Sol:-



player B

$m \times 2$

		B <sub>1</sub>	B <sub>2</sub>		
		A <sub>2</sub>	A <sub>3</sub>	5	5/9
		A <sub>3</sub>	A <sub>4</sub>	4	4/9
player A				3	6
				3/9	6/9



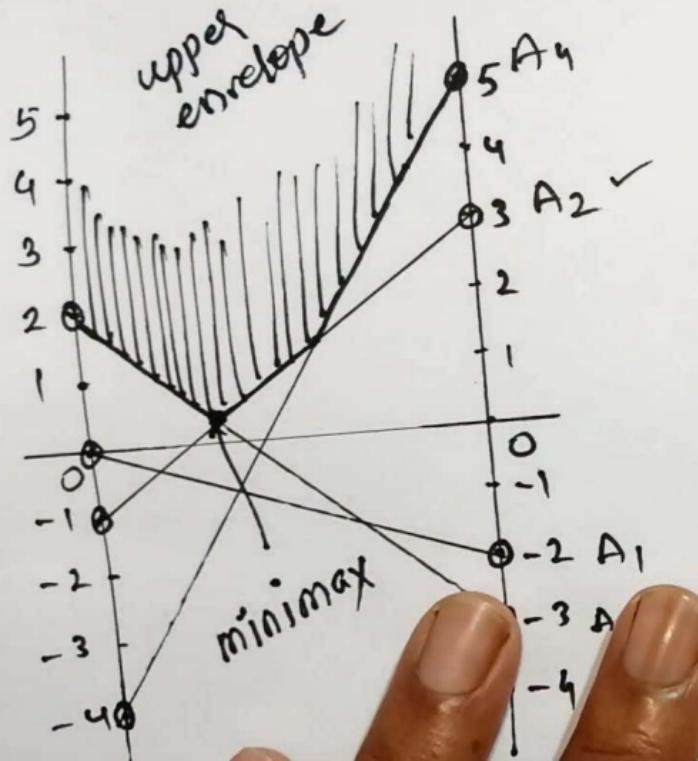
value of the game

$$\begin{aligned}
 v &= \beta \times \frac{5}{9} + (-\beta) \times \frac{4}{9} \\
 &= \frac{5}{3} - \frac{4}{3} \\
 &= \frac{1}{3}
 \end{aligned}$$

player A

$A_1$	-2		
$A_2$	3	-1	✓
$A_3$	-3	2	✓
$A_4$	5	-4	

$$\rightarrow A_3 \begin{bmatrix} -3 \\ 3 \\ 6 \\ 3/9 \end{bmatrix}$$



value of the game

$$\begin{aligned}
 v &= \beta \times \frac{5}{9} + (-\beta) \times \frac{4}{9} \\
 &= \frac{5}{3} - \frac{4}{3} \\
 &= \frac{1}{3}
 \end{aligned}$$

optimal strategy of player A  
 $= (0, 5/9, 4/9, 0)$

player B

$$= (3/9, 6/9)$$

## Game theory

### Simplex method - [m x n games]

Ques: Solve by Simplex method

$$A \begin{array}{c} B \\ \left[ \begin{array}{ccc} 3 & -2 & 4 \\ -1 & 4 & 2 \\ 2 & 2 & 6 \end{array} \right] \end{array}$$

Sol: Given matrix

$$A \begin{array}{c} B \\ \left[ \begin{array}{ccc} 3 & -2 & 4 \\ -1 & 4 & 2 \\ 2 & 2 & 6 \end{array} \right] \end{array}$$

② maximin

minimax

$\boxed{2-3}$

$$A \begin{bmatrix} 3 & -2 & 4 \\ -1 & 4 & 2 \\ 2 & 2 & 6 \end{bmatrix}$$

Sandeep Kumar Gour

Sol:: Given matrix

$$B$$

$$A \begin{bmatrix} 3 & -2 & 4 \\ -1 & 4 & 2 \\ 2 & 2 & 6 \end{bmatrix}$$

② maximin      ③ minimax

$v < 0$   
 $v = 0$   
 $v > 0$   
 $\boxed{2-3}$   
constant

We can infer that  $2 \leq v \leq 3$ . Hence it can be concluded that the value of the game lies between 2 and 3 and  $v > 0$

## Simplex method - [m x n games]

LPP

Ques: Solve by Simplex method

$$A \begin{bmatrix} 3 & -2 & 4 \\ -1 & 4 & 2 \\ 2 & 2 & 6 \end{bmatrix}$$

Sol: Given matrix

$$A \begin{bmatrix} x_1 & y_1 & y_2 & y_3 & -2 \\ x_1 & 3 & -2 & 4 & -2 \\ x_2 & -1 & 4 & 2 & -1 \\ x_3 & 2 & 2 & 6 & 2 \end{bmatrix} \quad \text{minmax}$$

gain

$$\begin{array}{l} V \rightarrow \min \\ V \rightarrow \max \\ \frac{x_1 + x_2 + x_3}{V} = \frac{1}{V} \\ \boxed{x_1 + x_2 + x_3 = \frac{1}{V}} \end{array}$$

$$\begin{array}{l} y_1 + y_2 + y_3 = 1 \\ y_1 + y_2 + y_3 = \frac{1}{V} \\ \boxed{y_1 + y_2 + y_3 = \frac{1}{V}} \end{array}$$

maximin

$$V > 0$$

2 - 3

const

$$\begin{array}{l} V < 0 \\ V = 0 \end{array}$$

We can infer that  $2 \leq V \leq 3$ . Hence it can be solved.

gain

$x_1$	-1	4	2	$\left[ \begin{array}{c} -1 \\ 4 \\ 2 \end{array} \right]$	② maximum
$x_2$	2	2	6		$v > 0$
③	4	6		$\boxed{2-3}$	<u>constant</u>

minimax

We can infer that  $2 \leq v \leq 3$ . Hence it can be concluded that the value of the game lies between 2 and 3 and  $v > 0$

LPP:  $\max \frac{1}{v} = y_1 + y_2 + y_3$

subject to constraints

$$3y_1 - 2y_2 + 4y_3 \leq 1$$

$$\frac{3y_1}{v} - 2\frac{y_2}{v} + 4\frac{y_3}{v} \leq \frac{1}{v}$$

minimax

2-3 constant

We can infer that  $2 \leq v \leq 3$ . Hence it can be concluded that the value of the game lies between 2 and 3 and  $v > 0$

LPP:  $\max \frac{1}{v} = y_1 + y_2 + y_3$

subject to constraints

$$3y_1 - 2y_2 + 4y_3 \leq 1$$

$$-1y_1 + 4y_2 + 2y_3 \leq 1$$

$$2y_1 + 2y_2 + 6y_3 \leq 1$$

$$\left| \begin{array}{l} \frac{3y_1}{v} - 2\frac{y_2}{v} + 4\frac{y_3}{v} \leq \frac{v}{v} \\ \end{array} \right.$$

and  $y_1, y_2, y_3 \geq 0$

$$\underline{\text{LPP}}: \max \frac{1}{\sqrt{v}} = y_1 + y_2 + y_3 \quad | \quad 3y_1 + 2y_2 + 4y_3 \leq v$$

$$\underline{\text{SLPP}}: \max \frac{1}{\sqrt{v}} = y_1 + y_2 + y_3 + 0s_1 + 0s_2 + 0s_3$$

s.t.

$$3y_1 - 2y_2 + 4y_3 + s_1 = 1$$

$$-1y_1 + 4y_2 + 2y_3 + s_2 = 1$$

$$2y_1 + 2y_2 + 6y_3 + s_3 = 1$$

$$\text{and } y_1, y_2, y_3, s_1, s_2, s_3 \geq 0$$

s.t.

$$3y_1 - 2y_2 + 4y_3 + s_1 = 1$$

$$-1y_1 + 4y_2 + 2y_3 + s_2 = 1$$

$$2y_1 + 2y_2 + 6y_3 + s_3 = 1$$

and  $y_1, y_2, y_3, s_1, s_2, s_3 \geq 0$

Now the following simplex table is formed

	$c_j$	1	1	1	0	0	0	Min Ratio
Basic variable	$c_B$	$y_B$	$y_1$	$y_2$	$y_3$	$s_1$	$s_2$	
$\leftarrow s_1$	0	1	③	-2	4	1	0	0
$s_2$	0	1	-1	4	2	0	1	0
$s_3$	0	1	2	2	6	0	0	1
	$\frac{1}{v} = 0$	$A_j$	-1↑	-1	-1	0	0	0
$y_1$	1	$y_3$	1	$-2/3$	$4/3$	$y_3$	0	0

Basic variable	$c_B$	$y_B$	$y_1$	$y_2$	$y_3$	$s_1$	$s_2$	$s_3$	Min Ratio
$\leftarrow s_1$	0	1	③	-2	4	1	0	0	$y_3 \leftarrow$
$s_2$	0	1	-1	4	2	0	1	0	-
$s_3$	0	1	2	2	6	0	0	1	$y_2$
$\frac{1}{v} = 0$		$\Delta_j$	-1↑	-1	-1	0	0	0	
$y_1$	1	$y_3$	1	$-2/3$	$4/3$	$y_3$	0	0	-
$s_2$	0	$4/3$	0	$10/3$	$10/3$	$y_3$	1	0	$2/5$
$\leftarrow s_3$	0	$1/3$	0	$10/3$	$10/3$	$-2/3$	0	1	$1/10 \leftarrow$
$\frac{1}{v} = \frac{1}{3}$		$\Delta_j$	0	$-5/3$ ↑	$y_3$	$y_3$	0	0	
$y_1$	1	$2/5$	1	0	2	$y_5$	0	$y_5$	
$s_2$	0	1	0	0	0	1	1	-1	
$y_2$	1	$1/10$	0	1	1	$-1/5$	0	$3/10$	
$\frac{1}{v} = \frac{1}{2}$		$\Delta_j$	0	0	2	0	0	$y_2$	

$\leftarrow s_3$	0	$1/\sqrt{3}$	0	$(10/3)$	$1/\sqrt{3}$	$-2/\sqrt{3}$	0	1	$\sqrt{10} \leftarrow$
$\frac{1}{\sqrt{v}} = \frac{1}{3}$	$\Delta j$	0	$-5/\sqrt{3} \uparrow$	$y_3$	$y_3$	0	0		
$y_1$	1	$2/\sqrt{5}$	1	0	2	$y_5$	0	$y_5$	
$s_2$	0	1	0	0	0	1	1	-1	
$y_2$	1	$1/\sqrt{10}$	0	1	1	$-1/\sqrt{5}$	0	$3/\sqrt{10}$	
$\frac{1}{\sqrt{v}} = \frac{1}{2}$	$\Delta j$	0	0	2	$0 \ 0 \ y_2$				

Thus, the solution for B's original problem is obtained as

$$y_1 = y_1 \times \sqrt{v} = \frac{2}{5} \times 2 = \frac{4}{5}$$

$$y_2 = y_2 \times \sqrt{v} = \frac{1}{\sqrt{10}} \times 2 = \frac{1}{\sqrt{5}}$$

$$y_3 = y_3 \times \sqrt{v} = 0 \times 2 = 0$$

$$y_1 = \frac{y_1}{\sqrt{v}}$$

$$y_1 = y_1 \times \sqrt{v}$$

$$\frac{1}{\sqrt{v}} = \frac{1}{2}$$

$$\text{i.e. } \boxed{v=2}$$

$$y_3 = Y_3 \times V = 0 \times 2 = 0$$

$$\frac{1}{V} = \frac{1}{2}$$

$$\therefore V = 2$$

The optimal strategies for the player A are obtained from the final table of the above problem. This is given by duality rules.

$$x_1 = X_1 \times V = 0 \times 2 = 0$$

$$x_1 = \frac{x_1}{V}$$

$$x_2 = X_2 \times V = 0 \times 2 = 0$$

$$x_2 = X_2 \times V$$

$$x_3 = X_3 \times V = \frac{1}{2} \times 2 = 1$$

The optimal strategies for the player A are obtained from the final table of the above problem. This is given by duality rules.

$$x_1 = x_1 \times \checkmark = 0 \times 2 = 0$$

$$x_2 = x_2 \times \checkmark = 0 \times 2 = 0$$

$$x_3 = x_3 \times \checkmark = \frac{1}{2} \times 2 = 1$$

$$x_1 = \frac{x_1}{\checkmark}$$

$$x_4 = x_1 \times \checkmark$$

therefore

optimal strategy of player A =  $(0, 0, 1)$

---

$$\text{B} = \left( \frac{4}{5}, \frac{1}{5}, 0 \right)$$

$$A \begin{bmatrix} 1 & -1 & -1 \\ -1 & -1 & 3 \\ -1 & 2 & -1 \end{bmatrix}$$

B

$$\boxed{v=0} \quad v > 0$$

$$\overbrace{-1 \quad +1}$$

Sol<sup>n</sup>: Given matrix

$$B \begin{bmatrix} 1 & -1 & -1 \\ -1 & -1 & 3 \\ -1 & 2 & -1 \end{bmatrix} \begin{matrix} -1 \\ -1 \\ -1 \end{matrix}$$

maximin = -1

$$\begin{matrix} 1 & 2 & 3 \end{matrix}$$

minimax

We can infer that  $-1 \leq v \leq 1$

since maximin value is -1, it is possible  
that value of the game may be negative

or

$$\begin{bmatrix} -1 & 2 & -1 \end{bmatrix}$$

$$\begin{array}{c} -1 \\ +1 \end{array}$$

Sol<sup>n</sup>: Given matrix

$$A = \begin{bmatrix} 1 & -1 & -1 \\ -1 & -1 & 3 \\ -1 & 2 & -1 \end{bmatrix} \xrightarrow{-1} \text{maximin} = -1$$

① 2 3  
minimax

We can infer that  $-1 \leq v \leq 1$

since maximin value is -1, it is possible  
that value of the game may be negative  
or zero, thus the constant c

$$\begin{bmatrix} 1 & 2 & -1 \end{bmatrix}$$

Sol<sup>n</sup>: Given matrix

$$A = \begin{bmatrix} 1 & -1 & -1 \\ -1 & -1 & 3 \\ -1 & 2 & -1 \end{bmatrix} \rightarrow \begin{array}{c} -1 \\ -1 \\ -1 \end{array}$$

maximin = -1

① 2 3

minimax

We can infer that  $-1 \leq v \leq 1$

since maximin value is -1, it is possible  
that value of the game may be negative  
or zero, thus the constant c is added to  
all the elements of matrix which is at  
least equal to the negative of maximin

$$A \begin{bmatrix} 1 & -1 & -1 \\ -1 & -1 & 3 \\ -1 & 2 & -1 \end{bmatrix} \begin{matrix} -1 \\ -1 \\ -1 \end{matrix} \text{ maximin} = -1$$

① 2 3

minimax

$\textcircled{c=1}, 2, \dots$

we can infer that  $-1 \leq v \leq 1$

since maximin value is -1, it is possible  
that value of the game may be negative  
or zero, thus the constant 'c' is added to  
all the elements of matrix which is at  
least equal to the negative of maximin.

let  $c=1$ , add this value to all the  
elements of the matrix.

The

$$A \begin{bmatrix} -1 & -1 & 3 \\ -1 & 2 & -1 \end{bmatrix}$$

$\checkmark = 0$

$$\begin{array}{c} -1 \\ +1 \end{array}$$

Sol<sup>n</sup>: Given matrix B

$$A \begin{bmatrix} 1 & -1 & -1 \\ -1 & -1 & 3 \\ -1 & 2 & -1 \end{bmatrix} \begin{matrix} -1 \\ -1 \\ -1 \end{matrix}$$

$$\text{maximin} = -1$$

$$\textcircled{1} \quad 2 \quad 3$$

minimax

$$\begin{array}{c} + \\ \textcircled{c=1}, 2, \dots \end{array}$$

B

$$A \begin{bmatrix} 2 & 0 & 0 \\ 0 & 0 & 4 \\ 0 & 3 & 0 \end{bmatrix}$$

$$A \begin{bmatrix} 2 & 0 & 0 \\ 0 & 0 & 4 \\ 0 & 3 & 0 \end{bmatrix} \quad B$$

$$Y_1 = \frac{y_1}{\sqrt{}}$$

LPP -  $\max \frac{1}{\sqrt{}} = Y_1 + Y_2 + Y_3$

subject to constraints

$$2Y_1 + 0Y_2 + 0Y_3 \leq 1$$

$$0Y_1 + 0Y_2 + 4Y_3 \leq 1$$

$$0Y_1 + 3Y_2 + 0Y_3 \leq 1$$

and  $Y_1, Y_2, Y_3 \geq 0$

	$C_B$	$C_J$	1	1	1	0	0	0	
Basic Variable	$C_B$	$C_J$	$Y_1$	$Y_2$	$Y_3$	$S_1$	$S_2$	$S_3$	Min Ratio $Y_B/Y_K$
$\leftarrow S_1$	0	1	(2)	0	0	1	0	0	$1/2 \leftarrow$
$S_2$	0	1	0	0	4	0	1	0	-
$S_3$	0	1	0	3	0	0	0	1	-
		$\frac{1}{\sqrt{v}} = 0$	$\Delta_j$	$-1 \uparrow$	-1	-1	0	0	
$Y_1$	1	$1/2$	1	0	0	$1/2$	0	0	-
$S_2$	0	1	0	0	4	0	1	0	-
$\leftarrow S_3$	0	1	0	(3)	0	0	0	1	$Y_3 \leftarrow$
		$\frac{1}{\sqrt{v}} = Y_2$	$\Delta_j$	0	-1	-1	$Y_2$	0	
$Y_1$	1	$1/2$	1	0	0	$1/2$	0	0	-
$\leftarrow S_2$	0	1	0	0	(4)	0	1	0	$1/4 \leftarrow$
$Y_2$	1	$1/3$	0	1	0	0	0	$Y_3$	-
		$\frac{1}{\sqrt{v}} = \frac{5}{6}$	$\Delta_j$	0	0	-1	$Y_2$	0	$Y_3$
$Y_1$	1	$1/2$	1	0	0	$1/2$	0	0	
$Y_2$	.	.	.	.	.	.	.	.	

$s_3$	0	1	0	3	0	0	0	1	-
	$\frac{1}{\sqrt{v}} = 0$	$\Delta j$	$-1 \uparrow$	-1	-1	0	0	0	
$y_1$	1	$1/2$	1	0	0	$1/2$	0	0	-
$s_2$	0	1	0	0	4	0	1	0	-
$\leftarrow s_3$	0	1	0	(3)	0	0	0	1	$y_3 \leftarrow$
	$\frac{1}{\sqrt{v}} = y_2$	$\Delta j$	0	-1	-1	$y_2$	0	0	
$y_1$	1	$1/2$	1	0	0	$1/2$	0	0	-
$\leftarrow s_2$	0	1	0	0	(4)	0	1	0	$1/4 \leftarrow$
$y_2$	1	$1/3$	0	1	0	0	0	$y_3$	-
	$\frac{1}{\sqrt{v}} = \frac{5}{6}$	$\Delta j$	0	0	-1	$1/2$	0	$1/3$	
$y_1$	1	$1/2$	1	0	0	$1/2$	0	0	
$y_3$	1	$1/4$	0	0	1	0	$1/4$	0	
$y_2$	1	$1/3$	0	1	0	0	0	$1/3$	
	$\frac{1}{\sqrt{v}} = \frac{13}{12}$	$\Delta j$	0	0	0	$y_2$	$1/4$	$1/3$	

$y_1$	1	$\frac{1}{2}$	1	0	0	$\frac{1}{2}$	0	0	-
$s_2$	0	1	0	0	4	0	1	0	$\frac{1}{4} \leftarrow$
$y_2$	1	$\frac{1}{3}$	0	1	0	0	0	$y_3$	-
	$\frac{1}{v} = \frac{5}{6}$	$\Delta j$	0	0	-1	$\frac{1}{2}$	0	$y_3$	
$y_1$	1	$\frac{1}{2}$	1	0	0	$\frac{1}{2}$	0	0	
$y_3$	1	$\frac{1}{4}$	0	0	1	0	$\frac{1}{4}$	0	
$y_2$	1	$\frac{1}{3}$	0	1	0	0	0	$y_3$	
	$\frac{1}{v} = \frac{13}{12}$	$\Delta j$	0	0	0	$y_2$	$\frac{1}{4}$	$y_3$	

$$\text{Here } \frac{1}{v} = \frac{13}{12}$$

$\Delta u \quad \Delta g \quad \Delta e$

$$\text{i.e. } v = 12/13$$

$$y_1 = y_1 \times v = \frac{1}{2} \times \frac{12}{13} = \frac{6}{13}$$

$$y_2 = y_2 \times v = \frac{1}{3} \times \frac{12}{13} = \frac{4}{13}$$

$$y_3 = y_3 \times v = \frac{1}{4} \times \frac{12}{13} = \frac{3}{13}$$

$$x_1 = x_1 \times v = \frac{1}{2} \times \frac{12}{13} = \frac{6}{13}$$

$$x_2 = x_2 \times v = \frac{1}{4} \times \frac{12}{13} = \frac{3}{13}$$

$$x_3 = x_3 \times v = \frac{1}{3} \times \frac{12}{13} = \frac{4}{13}$$

$y_1$							$\frac{1}{4}$	0
$y_3$	1	$\frac{1}{4}$	0	0	1	0	$\frac{1}{4}$	0
$y_2$	1	$\frac{1}{3}$	0	1	0	0	0	$\frac{1}{3}$
	$\frac{1}{v} = \frac{13}{12}$	$\Delta_j$	0	0	0	$y_2$	$\frac{1}{4}$	$\frac{1}{3}$

Here  $\frac{1}{v} = \frac{13}{12}$

$\Delta_u$   $\Delta_g$   $\Delta_e$

i.e  $v = 12/13$

$$\left. \begin{array}{l} y_1 = y_1 \times v = \frac{1}{2} \times \frac{12}{13} = \frac{6}{13} \\ y_2 = y_2 \times v = \frac{1}{3} \times \frac{12}{13} = \frac{4}{13} \\ \vdots \quad \vdots \quad \vdots \end{array} \right| \quad \left. \begin{array}{l} x_1 = x_1 \times v = \frac{1}{2} \times \frac{12}{13} = \frac{6}{13} \\ x_2 = x_2 \times v = \frac{1}{4} \times \frac{12}{13} = \frac{3}{13} \end{array} \right.$$

optimal strategy of player A =  $(\frac{6}{13}, \frac{3}{13}, \frac{4}{13})$

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B =  $(\frac{6}{13}, \frac{4}{13}, \frac{3}{13})$

value of game =  $v - c = \frac{12}{13} - 1$   
 $= -\frac{1}{13}$       Ans.