

## Operations Research -

operations Research is the systematic, method-oriented study of the basic structure, characteristics, functions and relationships of an organization to provide the executive with a sound, scientific and quantitative basis for decision making.

## Linear Programming Problem (LPP) -

A linear form is meant a mathematical expression of the type  $a_1x_1 + a_2x_2 + \dots + a_nx_n$ , where  $a_1, a_2, \dots, a_n$  are constants and  $x_1, x_2, \dots, x_n$  are variables. The term 'programming' refers to the process of determining a particular programme or plan of action.

Linear programming deals with the optimization (Maximization/Minimization) of a function of variables known as objective function, subject to a set of linear equations and/or inequalities known as constraints or restrictions.

## General Form of Linear Programming Problem (LPP) -

**Definition:-** The general LPP calls for optimizing (Maximizing or minimizing) a linear function of variables called the '**Objective function**' subject to a set of linear equations and/or inequalities called the '**constraints**' or '**Restrictions**'.

**General Form:**

Objective Function (Max or Min)  $Z = c_1x_1 + c_2x_2 + \dots + c_nx_n$

subject to constraints  $a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n (\leq or \geq) b_1$   
 $a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n (\leq or \geq) b_2$   
 $\vdots$   
 $a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n (\leq or \geq) b_m$

and  $x_1, x_2, \dots, x_n \geq 0$

## Formulation of LP problems -

Eg: A firm manufactures two types of products A and B sells them at a profit of Rs. 2 on type A and Rs. 3 on type B. Each product is processed on two machines G and H. Type A requires 1 minute of processing time on G and 2 minutes on H; type B requires 1 minute on G and 1 minute on H. The machine G is available for not more than 6 hours 40 minutes while machine H is available for 10 hours during any working day. Formulate the problem as LPP.

Sol<sup>n</sup>:  $x_1$  be the number of products of type A  
 $x_2$  be the number of products of type B

Machine	Type of products (minutes)		
	Type A ( $x_1$ units)	Type B ( $x_2$ units)	Available time (minutes)
G	1	1	400
H	2	1	600
profit per unit	Rs. 2	Rs. 3	

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profit per unit	Rs. 2	Rs. 3	

$$2x_1 \quad 3x_2$$

$$2x_1 + 3x_2 \quad (\text{objective function})$$

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profit per unit	Rs. 2	Rs. 3	

$$2x_1$$

$$3x_2$$

$$\text{Max } Z = 2x_1 + 3x_2 \quad (\text{objective function})$$

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Profit per unit	Rs. 2	Rs. 3	

$$2x_1 \quad 3x_2$$

$$\text{Max } Z = 2x_1 + 3x_2 \quad (\text{objective function})$$

$$\begin{cases} x_1 + x_2 \leq 400 \\ 2x_1 + x_2 \leq 600 \end{cases} \quad \left. \begin{array}{l} \text{constraints} \\ \text{non-negative constraints} \end{array} \right\}$$

non-negative constraints are

$$x_1 \geq 0, x_2 \geq 0$$

Profit per unit	Rs. 2	Rs. 3	
-----------------	-------	-------	--

$$2x_1$$

$$3x_2$$

$$\text{Max } Z = 2x_1 + 3x_2 \quad (\text{objective function})$$

$$\begin{cases} x_1 + x_2 \leq 400 \\ 2x_1 + x_2 \leq 600 \end{cases} \quad \text{constraints}$$

non-negative constraints are

$$x_1 \geq 0, x_2 \geq 0$$

LPP

$$\text{Max } Z = 2x_1 + 3x_2$$

subject to constraints

$$x_1 + x_2 \leq 400$$

$$2x_1 + x_2 \leq 600$$

$$\text{and } x_1, x_2 \geq 0$$

## Graphical Solution of LP problems:-

Q. Max  $Z = 3x_1 + 4x_2$

Subject to constraints

$$x_1 + x_2 \leq 450 \quad (1)$$

$$2x_1 + x_2 \leq 600 \quad (2)$$

and  $x_1, x_2 \geq 0$

Soln: from constraints (1)

$$x_1 + x_2 = 450$$

$$\begin{aligned} \text{put } x_1 = 0 &\Rightarrow 0 + x_2 = 450 \\ &\Rightarrow x_2 = 450 \end{aligned}$$

$$\text{put } x_2 = 0 \Rightarrow x_1 = 450$$

from constraints (2)

$$2x_1 + x_2 = 600$$

$$\begin{aligned} \text{put } x_1 = 0 &\Rightarrow 0 + x_2 = 600 \\ &\Rightarrow x_2 = 600 \end{aligned}$$

$$\begin{aligned} \text{put } x_2 = 0 &\Rightarrow 2x_1 = 600 \\ &\Rightarrow x_1 = 300 \end{aligned}$$

## solution of LP problems:-

Q. Max  $Z = 3x_1 + 4x_2$

Subject to constraints

$$x_1 + x_2 \leq 450 \quad (1)$$

$$2x_1 + x_2 \leq 600 \quad (2)$$

and  $x_1, x_2 \geq 0$

Soln: from constraints (1)

$$x_1 + x_2 = 450$$

$$\begin{aligned} \text{put } x_1 = 0 &\Rightarrow 0 + x_2 = 450 \\ &\Rightarrow x_2 = 450 \end{aligned}$$

$$\begin{aligned} \text{put } x_2 = 0 &\Rightarrow x_1 = 450 \\ \text{from constraints (2)} & \end{aligned}$$

$$2x_1 + x_2 = 600$$

$$\begin{aligned} \text{put } x_1 = 0 &\Rightarrow 0 + x_2 = 600 \\ &\Rightarrow x_2 = 600 \end{aligned}$$

$$\begin{aligned} \text{put } x_2 = 0 &\Rightarrow 2x_1 = 600 \\ &\Rightarrow x_1 = 300 \end{aligned}$$

constraints (1)

$x_1$	$x_2$
0	450
450	0

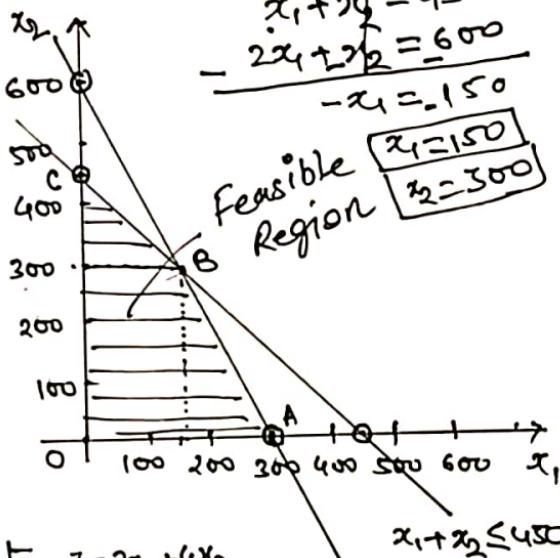
constraints (2)

$x_1$	$x_2$
0	600
300	0

$$\begin{aligned} x_1 + x_2 &= 450 \\ -2x_1 + x_2 &= 600 \\ -x_1 &= -150 \end{aligned}$$

$$\begin{aligned} x_1 &= 150 \\ x_2 &= 300 \end{aligned}$$

feasible Region



points coordinate  $Z = 3x_1 + 4x_2$

$$O (0, 0) = 0$$

$$A (300, 0) = 900$$

$$B (150, 300) = 1650$$

$$C (0, 450) = 1800$$

$$\begin{aligned} x_1 + x_2 &\leq 450 \\ 2x_1 + x_2 &\leq 600 \end{aligned}$$

optimal solution Max  $Z = 1800$  at  $x_1 = 0$   
 $x_2 = 450$

## Graphical Solution using Objective Function Line :-

$$\text{Min } Z = 20x_1 + 10x_2$$

Subject to constraints

$$x_1 + 2x_2 \leq 40 \quad (1)$$

$$3x_1 + x_2 \geq 30 \quad (2)$$

$$4x_1 + 3x_2 \geq 60 \quad (3)$$

$$\text{and } x_1, x_2 \geq 0$$

Sol<sup>n</sup>: The first constraint written  
in form of equation

$$x_1 + 2x_2 = 40$$

$$\begin{aligned} \text{put } x_1 = 0 &\Rightarrow 0 + 2x_2 = 40 \\ &\Rightarrow x_2 = 20 \end{aligned}$$

$$\begin{aligned} \text{put } x_2 = 0 &\Rightarrow x_1 + 0 = 40 \\ &\Rightarrow x_1 = 40 \end{aligned}$$

The second constraint written  
in form of equation

$$3x_1 + x_2 = 30$$

$$\begin{aligned} \text{put } x_1 = 0 &\Rightarrow 0 + x_2 = 30 \\ &\Rightarrow x_2 = 30 \end{aligned}$$

$$\begin{aligned} \text{put } x_2 = 0 &\Rightarrow 3x_1 + 0 = 30 \\ &\Rightarrow x_1 = 10 \end{aligned}$$

The third constraint written  
in form of equation

$$4x_1 + 3x_2 = 60$$

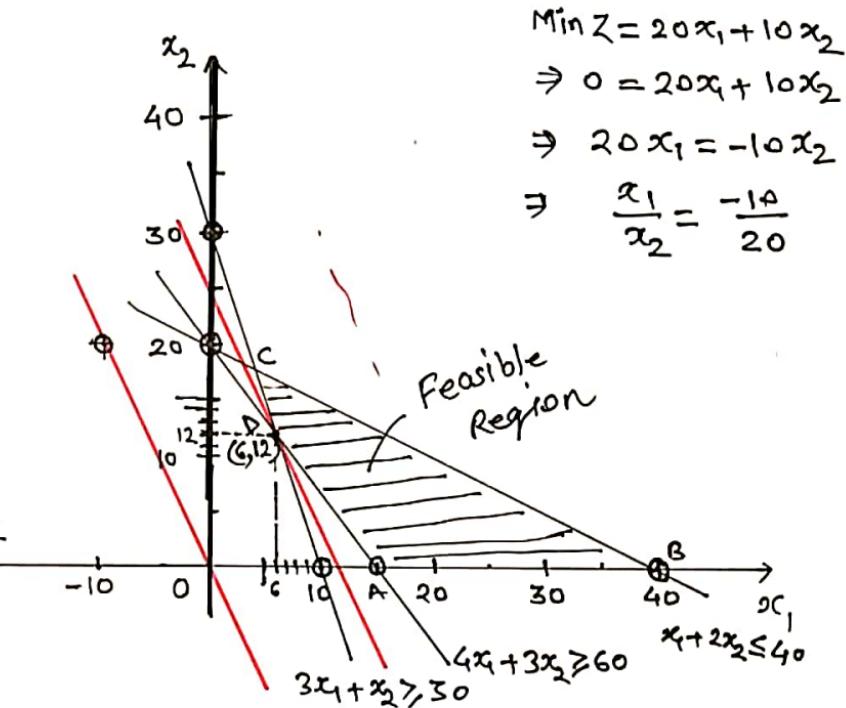
$$\begin{aligned} \text{put } x_1 = 0 &\Rightarrow 0 + 3x_2 = 60 \\ &\Rightarrow x_2 = 20 \end{aligned}$$

$$\begin{aligned} \text{put } x_2 = 0 &\Rightarrow 4x_1 + 0 = 60 \\ &\Rightarrow x_1 = 15 \end{aligned}$$

Constraints	$x_1$	$x_2$
1. $x_1 + 2x_2 \leq 40$	0	20
	40	0
2. $3x_1 + x_2 \geq 30$	0	30
	10	0
3. $4x_1 + 3x_2 \geq 60$	0	20
	15	0

optimal solution

$$\begin{aligned} \text{Min } Z &= 20x_1 + 10x_2 \\ &= 120 + 120 \\ &= \underline{\underline{240}} \end{aligned}$$

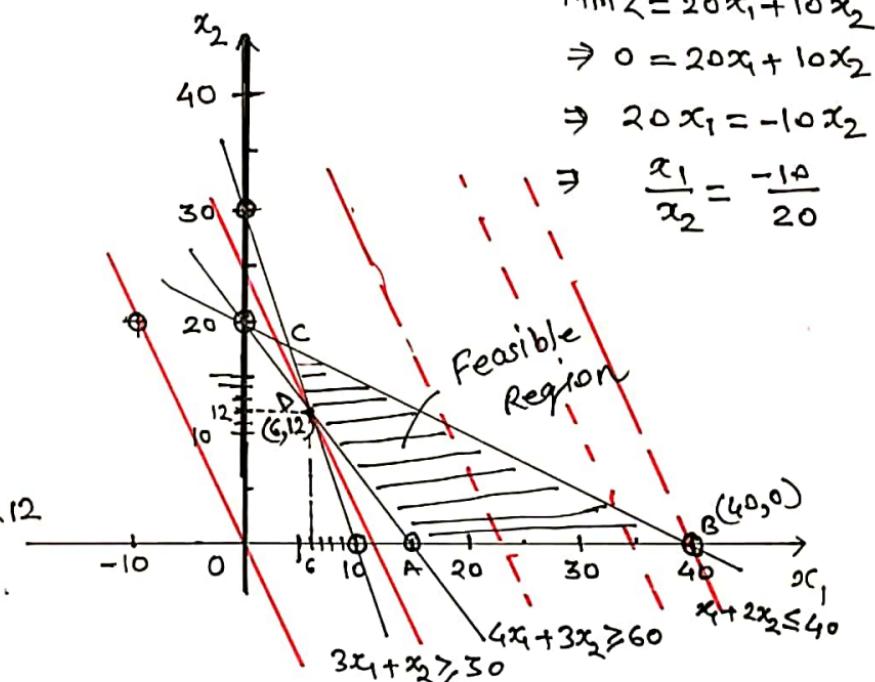


Constraints	$x_1$	$x_2$
1. $x_1 + 2x_2 \leq 40$	0	20
	40	0
2. $3x_1 + x_2 \geq 30$	0	30
	10	0
3. $4x_1 + 3x_2 \geq 60$	0	20
	15	0

optimal solution

$$\begin{aligned}
 \text{Min } Z &= 20x_1 + 10x_2 \\
 &= 120 + 120 \\
 &= \underline{\underline{240}}
 \end{aligned}$$

$$\begin{aligned}
 \text{Max } Z &= 20x_1 + 10x_2 \\
 &= 800
 \end{aligned}$$



$$\begin{aligned}
 \text{Min } Z &= 20x_1 + 10x_2 \\
 \Rightarrow 0 &= 20x_1 + 10x_2 \\
 \Rightarrow 20x_1 &= -10x_2 \\
 \Rightarrow \frac{x_1}{x_2} &= \frac{-10}{20}
 \end{aligned}$$

## Multiple Optimal Solution in Graphical Method -

$$\text{Max } Z = 4x_1 + 3x_2$$

Subject to constraints

$$4x_1 + 3x_2 \leq 24 \quad (1)$$

$$x_1 \leq 4.5 \quad (2)$$

$$x_2 \leq 6 \quad (3)$$

$$\text{and } x_1, x_2 \geq 0$$

Sol: The first constraint written  
in a form of equation

$$4x_1 + 3x_2 = 24$$

$$\text{put } x_1=0 \Rightarrow 0 + 3x_2 = 24$$

$$\Rightarrow x_2 = 8$$

$$\text{put } x_2=0 \Rightarrow 4x_1 + 0 = 24$$

$$\Rightarrow x_1 = 6$$

The second constraint written  
in a form of equation

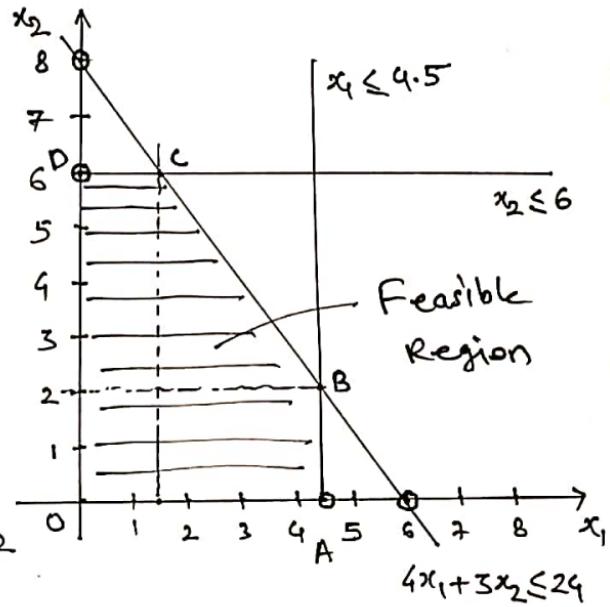
$$x_1 = 4.5, x_2 = 0$$

The third constraint written  
in a form of equation

$$x_2 = 6, x_1 = 0$$

constraints	$x_1$	$x_2$
1. $4x_1 + 3x_2 \leq 24$	0	8
	6	0
2. $x_1 \leq 4.5$	4.5	0
3. $x_2 \leq 6$	0	6

constraints	$x_1$	$x_2$
1. $4x_1 + 3x_2 \leq 24$	0	8
	6	0
2. $x_1 \leq 4.5$	4.5	0
3. $x_2 \leq 6$	0	6



points coordinate  $\text{Max } Z = 4x_1 + 3x_2$

0	(0, 0)	= 0
A	(4.5, 0)	= 18
B	(4.5, 2)	= 24
C	(1.5, 6)	= 24
D	(0, 6)	= 18

optimal solution  
 $\text{Max } Z = 24$   
at point B(4.5, 2)  
point C(1.5, 6).

## No optimal solution in Graphical Method -

Qn:  $\text{Max } Z = 3x_1 + 2x_2$

Subject to constraints

$$x_1 + x_2 \leq 1 \quad (1)$$

$$x_1 + x_2 \geq 3 \quad (2)$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: from constraint (1)

$$x_1 + x_2 = 1$$

$$\text{put } x_1 = 0 \Rightarrow x_2 = 1 \quad (0, 1)$$

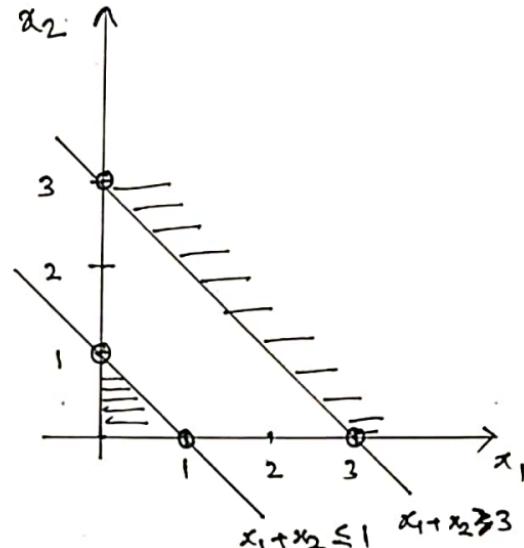
$$\text{put } x_2 = 0 \Rightarrow x_1 = 1 \quad (1, 0)$$

from constraint (2)

$$x_1 + x_2 = 3$$

$$\text{put } x_1 = 0 \Rightarrow x_2 = 3 \quad (0, 3)$$

$$\text{put } x_2 = 0 \Rightarrow x_1 = 3 \quad (3, 0)$$



## No optimal solution in Graphical Method -

Qn:  $\text{Max } Z = 3x_1 + 2x_2$

Subject to constraints

$$x_1 + x_2 \leq 1 \quad (1)$$

$$x_1 + x_2 \geq 3 \quad (2)$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: from constraint (1)

$$x_1 + x_2 = 1$$

$$\text{put } x_1 = 0 \Rightarrow x_2 = 1 \quad (0, 1)$$

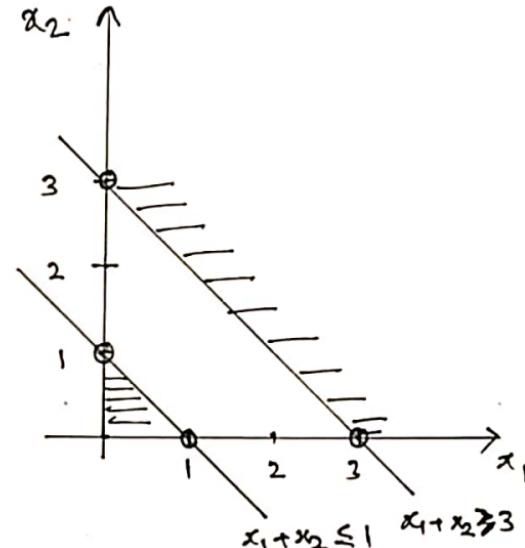
$$\text{put } x_2 = 0 \Rightarrow x_1 = 1 \quad (1, 0)$$

from constraint (2)

$$x_1 + x_2 = 3$$

$$\text{put } x_1 = 0 \Rightarrow x_2 = 3 \quad (0, 3)$$

$$\text{put } x_2 = 0 \Rightarrow x_1 = 3 \quad (3, 0)$$



There is no feasible region.

Hence there is 'No optimal' solution.

## Unbounded solution in Graphical Method -

Q: Max  $Z = 3x_1 + 2x_2$

Subject to constraints

$$x_1 - x_2 \leq 1 \quad \text{--- (1)}$$

$$x_1 + x_2 \geq 3 \quad \text{--- (2)}$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: from constraint (1)

$$x_1 - x_2 = 1$$

put  $x_1 = 0 \Rightarrow x_2 = -1 \quad (0, -1)$

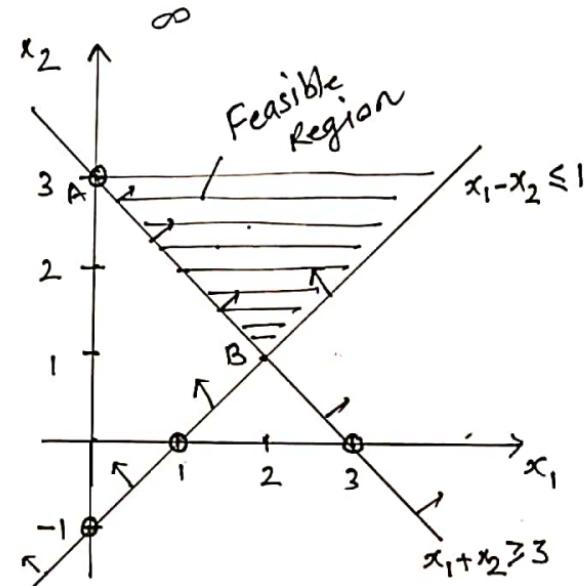
put  $x_2 = 0 \Rightarrow x_1 = 1 \quad (1, 0)$

from constraint (2)

$$x_1 + x_2 = 3$$

put  $x_1 = 0 \Rightarrow x_2 = 3 \quad (0, 3)$

put  $x_2 = 0 \Rightarrow x_1 = 3 \quad (3, 0)$



The maximum value of the objective function occurs at a point at  $\infty$ . Hence the given problem has Unbounded solution.

## General LPP to Standard LPP -

1. Write the objective function in the maximization form.
2. Convert all inequalities as equations.
3. The right side element of each constraint should be made non-negative.
4. All variables must have non-negative values.

## General LPP to Standard LPP -

1. Write the objective function in the maximization form.

$$\text{Max } Z \quad \text{Min } Z \quad [\text{Max } Z' = \text{Min}(-Z)] \quad \text{Eg: } \text{Min } Z = -3x_1 + x_2 \\ \text{Max } Z' = 3x_1 - x_2$$

2. Convert all inequalities as equations.

$$\leq \text{ + slack} \quad \left| \begin{array}{l} \text{Eg: } x_1 + 2x_2 \leq 12 \quad \& \quad 2x_1 + x_2 \geq 15 \\ x_1 + 2x_2 + s_1 = 12 \quad \quad \quad 2x_1 + x_2 - s_2 = 15 \end{array} \right. \\ \geq \text{ - surplus}$$

3. The right side element of each constraint should be made non-negative

$$\text{Eg: } 2x_1 + x_2 - s_2 = -15 \quad (\text{Multiplying by } -1) \\ \Rightarrow -2x_1 - x_2 + s_2 = 15$$

4. All variables must have non-negative values.

$$\text{Eg: } x_1 + x_2 \leq 3, \quad x_1 \geq 0, \quad x_2 \text{ is } \underline{\text{unrestricted}} \text{ in sign}$$

$$x_1 + (x_2' - x_2'') \leq 3 \quad x_1, x_2', x_2'' \geq 0$$

## General LPP to Standard LPP -

1. Write the objective function in the maximization form.

$$\text{Max } Z \quad \text{Min } Z \quad [\text{Max } Z' = \text{Min}(-Z)] \quad \text{Eg: } \text{Min } Z = -3x_1 + x_2 \\ \text{Max } Z' = 3x_1 - x_2$$

2. Convert all inequalities as equations.

$$\begin{array}{ll|l} \leq & + \text{slack} & \text{Eg: } x_1 + 2x_2 \leq 12 \quad \& \quad 2x_1 + x_2 \geq 15 \\ \geq & - \text{surplus} & x_1 + 2x_2 + s_1 = 12 \quad \quad \quad 2x_1 + x_2 - s_2 = 15 \end{array}$$

3. The right side element of each constraint should be made non-negative

$$\text{Eg: } 2x_1 + x_2 - s_2 = -15 \quad (\text{Multiplying by } -1) \\ \Rightarrow -2x_1 - x_2 + s_2 = 15$$

4. All variables must have non-negative values.

$$\text{Eg: } x_1 + x_2 \leq 3, \quad x_1 \geq 0, \quad x_2 \text{ is } \underline{\text{unrestricted}} \text{ in sign}$$

$$x_1 + (x_2' - x_2'') \leq 3 \quad x_1, x_2', x_2'' \geq 0$$

$$x_1 + x_2' - x_2'' + s_1 = 3$$

## Examples: GLPP to SLPP -

$$\textcircled{1} \quad \text{Max } Z = 3x_1 + x_2$$

Subject to

$$2x_1 + x_2 + s_1 \leq 2$$

$$3x_1 + 4x_2 \geq 12$$

$$\text{and } x_1, x_2 \geq 0$$

### SLPP

$$\text{Max } Z = 3x_1 + x_2$$

s.t.

$$2x_1 + x_2 + s_1 = 2$$

$$3x_1 + 4x_2 - s_2 = 12$$

$$\text{and } x_1, x_2, s_1, s_2 \geq 0$$

$$\textcircled{2} \quad \text{Min } Z = 4x_1 + 2x_2$$

subject to

$$3x_1 + x_2 \geq 2$$

$$x_1 + x_2 \geq 21$$

$$x_1 + 2x_2 \geq 30$$

$$\text{and } x_1, x_2 \geq 0$$

### SLPP

$$\text{Max } Z' = -4x_1 - 2x_2$$

subject to

$$3x_1 + x_2 - s_1 = 2$$

$$x_1 + x_2 - s_2 = 21$$

$$x_1 + 2x_2 - s_3 = 30$$

$$\text{and } x_1, x_2, s_1, s_2, s_3 \geq 0$$

$$\textcircled{3} \quad \text{Min } Z = x_1 + 2x_2 + 3x_3$$

subject to

$$2x_1 + 3x_2 + 3x_3 \geq -4$$

$$3x_1 + 5x_2 + 2x_3 \leq 7$$

$$\text{and } x_1, x_2 \geq 0,$$

$x_3$  is unrestricted in sign

### SLPP

$$\text{Max } Z' = -x_1 - 2x_2 - 3x_3$$

subject to

$$-2x_1 - 3x_2 - 3(x_3^I - x_3^{II}) + s_1 = 4$$

$$3x_1 + 5x_2 + 2(x_3^I - x_3^{II}) + s_2 = 7$$

$$\text{and } x_1, x_2, x_3^I, x_3^{II}, s_1, s_2 \geq 0$$

Step 4: Calculation of  $Z$  and  $\Delta_j$  (Net Evaluation) and test the basic feasible solution for optimality by the rules given.

$$Z = C_B X_B \quad \text{and} \quad \Delta_j = Z_j - C_j = c_B x_j - c_j$$

Rule 1 - If all  $\Delta_j \geq 0$ , the solution under the test will be optimal.

Rule 2 - If at least one  $\Delta_j$  is negative, the solution is not optimal and then proceeds to improve the solution in the next step.

Rule 3 - If corresponding to any negative  $\Delta_j$ , all elements of the column  $x_j$  are negative or zero, then the solution under test will be unbounded.

### Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

$$\text{and } x_1, x_2 \geq 0$$

Sol<sup>n</sup>: Step① SLP

### Step② Construct starting simplex table

Basic Variable	$C_B$	$C_j$	80	55	0	0
$s_1$	0	$x_B$	40	4	2	1
$s_2$	0	$x_1$	32	2	4	0

$$Z = 0 \times 40 + 8 \times 32$$

$$= 0$$

$$\begin{aligned}\Delta_1 &= C_B x_1 - C_1 \\ &= (0,0)(4,2) - 80 \\ &= 0 + 0 - 80 = -80\end{aligned}$$

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## Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① SLPP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

$$4x_1 + 2x_2 + s_1 = 40$$

$$2x_1 + 4x_2 + s_2 = 32$$

and  $x_1, x_2, s_1, s_2 \geq 0$

Step② Represent in matrix form

$$AX = B$$

$$\begin{bmatrix} x_1 & x_2 & s_1 & s_2 \\ 4 & 2 & 1 & 0 \\ 2 & 4 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 40 \\ 32 \\ 0 \end{bmatrix}$$

Step③ Construct starting simplex table

Basic Variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$C_j$
$s_1$	0	40	4	2	1	0	80
$s_2$	0	32	2	4	0	1	55

Step 4: Calculation of  $Z$  and  $\Delta_j$  (Net Evaluation) and test the basic feasible solution for optimality by the rules given.

$$Z = C_B X_B \quad \text{and} \quad \Delta_j = Z_j - C_j = C_B X_j - C_j$$

### Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

$$\text{and } x_1, \quad Z = 0 \times 40 + 0 \times 32 \\ = 0$$

Sol<sup>n</sup>: Step①

$$\text{Max } Z =$$

Subject to

$$4x_1 +$$

$$2x_1 +$$

and

Step

Step⑤ construct starting simplex table

	$C_j$	80	55	0	0	
Basic Variable	$C_B$	$x_0$	$x_1$	$x_2$	$s_1$	$s_2$
$s_1$	0	40	4	2	1	0
$s_2$	0	32	2	4	0	1
$Z = 0$		$\Delta_j \rightarrow -80$	-55	0	0	

$$\Delta_1 = C_B x_1 - C_1$$

$$= (0,0)(4,2) - 80$$

$$= 0 + 0 - 80 = -80$$

$$\Delta_2 = C_B x_2 - C_2$$

$$= (0,0)(2,4) - 55$$

$$= 0 + 0 - 55 = -55$$

$$\Delta_3 = C_B x_3 - C_3$$

$$= (0,0)(1,0) - 0$$

$$= 0 + 0 - 0 = 0$$

$$\Delta_4 = C_B x_4 - C_4$$

$$= (0,0)(0,1) - 0$$

$$= 0 + 0 - 0 = 0$$

Step 5: To improve the basic feasible solution, the incoming and outgoing vectors are determined.

(i) Incoming (Entering) Vector:  $\overset{(x_k)}{\rightarrow}$  most negative value of  $A_{ij}^{(k)}$

(ii) Outgoing (Leaving) Vector:  $\min [x_0/x_k \rightarrow x_k > 0]$

### Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① SLPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subj. to

$$+ 2x_2$$

$$\leq 40$$

$$\leq 32$$

$$+ 6x_2$$

Step② construct starting simplex table

Basic Variable	$C_B$	$x_0$	$x_1$	$x_2$	$s_1$	$s_2$	$C_j$
$s_1$	0	40	4	2	1	0	$-80$
$s_2$	0	32	2	4	0	1	$-55$
$Z = 0$							0

Step 5: To improve the basic feasible solution, the incoming and outgoing vectors are determined.

- (i) Incoming (Entering) Vector:  $x_k^{(k)}$  most negative value of  $A_{ij}^{(k)}$
- (ii) Outgoing (Leaving) Vector:  $\min [x_B/x_k \rightarrow x_k > 0]$

### Simplex Method -

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$$\text{Max } Z = 80x_1 + 55x_2$$

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and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① SLPP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

$$4x_1 + 2x_2 + s_1 = 40$$

$$2x_1 + 4x_2 + s_2 = 32$$

Step② construct starting simplex table

Basic Variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$x_B/x_k$	$C_j - 80 \quad 55 \quad 0 \quad 0$	Min R.
$\leftarrow s_1$	0	40	4	2	1	0	$40/4=10$		
$s_2$	0	32	2	4	0	1	$32/2=16$		
$Z = 0$			$\Delta_j \rightarrow -80$	-55	0	0			

outgoing vector

( $x_k$ ) incoming vector.

Step 6: Mark the key element (pivot element) at the intersection of outgoing vector and incoming vector.

(i) Key element should be one.

(ii) Remaining values of  $x_k$  should be zero.  
(column)

Step 7: Now repeat step 4 to step 6 until an optimal solution is obtained.

### Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step ① SLP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

Step ② Construct starting simplex table

	$C_j$	80	55	0	0	MmR	
Basic Variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$x_B/x_k$
$\leftarrow s_1$	0	40	4	2	1	0	$40/4=10$
$s_2$	0	32	2	4	0	1	$32/2=16$
$Z = 0$		$\Delta j \rightarrow -80$	-55	0	0		

outgoing vector

( $x_2$ ) incoming vector.

(column)

Step 7: Now repeat step 4 to step 6 until an optimal solution is obtained.

### Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① SLPP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

$$4x_1 + 2x_2 + s_1 = 40$$

$$2x_1 + 4x_2 + s_2 = 32$$

and  $x_1, x_2, s_1, s_2 \geq 0$

Step 2) Convert into standard form

Step 3) construct starting simplex table

Basic Variable	C <sub>B</sub>	C <sub>j</sub>	80	55	0	0	MmR.
$\leftarrow s_1$	0	40	4	2	1	0	$40/4=10$
$s_2$	0	32	2	4	0	1	$32/2=16$
$Z = 0$		$\Delta_j \rightarrow$	-80	-55	0	0	

outgoing vector

( $x_2$ ) incoming vector.

Basic Variable	C <sub>B</sub>	C <sub>j</sub>	80	55	0	0
$x_1$	80	10	1	$1/2$	$1/4$	0
$s_2$	0	12	0	3	$-1/2$	1

## Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① SLPP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

$$4x_1 + 2x_2 + s_1 = 40$$

$$2x_1 + 4x_2 + s_2 = 32$$

and  $x_1, x_2, s_1, s_2 \geq 0$

Step② Represent in Matrix form

$$AX = B$$

$$\begin{bmatrix} x_1 & x_2 & s_1 & s_2 \\ 4 & 2 & 1 & 0 \\ 2 & 4 & 0 & 1 \end{bmatrix} \begin{bmatrix} 40 \\ 32 \\ 0 \end{bmatrix}$$

Step③ construct starting simplex table

	$C_j$	80	55	0	0	MinR.	
Basic Variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$x_B/x_R$
$\leftarrow s_1$	0	40	4	2	1	0	$40/4=10$
$s_2$	0	32	2	4	0	1	$32/2=16$
$Z = 0$		$\Delta j \rightarrow -80$	-55	0	0		

outgoing vector

( $x_2$ ) incoming vector

	$C_j$	80	55	0	0	minR.	
Basic Variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$x_B/x_R$
$x_1$	80	10	1	$\frac{1}{2}$	$\frac{1}{4}$	0	20
$\leftarrow s_2$	0	12	0	3	$-\frac{1}{2}$	1	4 ←
$Z = 800$		$\Delta j$	0	$+15\frac{1}{2}$	20	0	incoming

outgoing vector

$$2x_1 + 4x_2 \leq 32$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① SLPP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

$$4x_1 + 2x_2 + s_1 = 40$$

$$2x_1 + 4x_2 + s_2 = 32$$

and  $x_1, x_2, s_1, s_2 \geq 0$

Step② Represent in Matrix form

$$AX = B$$

$$\begin{bmatrix} x_1 & x_2 & s_1 & s_2 \\ 4 & 2 & 1 & 0 \\ 2 & 4 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 40 \\ 32 \end{bmatrix}$$

(x<sub>2</sub>) incoming vector.

Basic Variable	C <sub>B</sub>	C <sub>1</sub>	C <sub>2</sub>	C <sub>3</sub>	C <sub>4</sub>	C <sub>5</sub>	minR
$x_1$	80	10	1	$\frac{1}{2}$	$\frac{1}{4}$	0	20
$s_2$	0	12	0	3	$-\frac{1}{2}$	1	4
Z = 800		A <sub>j</sub>	0	-15	20	0	incoming

Basic Variable	C <sub>B</sub>	C <sub>1</sub>	C <sub>2</sub>	C <sub>3</sub>	C <sub>4</sub>	C <sub>5</sub>
$x_1$	80					
$x_2$	55	4	0	1	$-\frac{1}{6}$	$\frac{1}{3}$

$$R_1 = R_1 - \frac{1}{2}R_2$$

$$\frac{1}{4} + \frac{1}{12}$$

$$\frac{3+1}{12} = \frac{4}{12}$$

$$\begin{array}{r}
\begin{array}{rrrrr}
10 & 1 & \frac{1}{2} & \frac{1}{4} & 0 \\
- & 2 & 0 & \frac{1}{2} & + \frac{1}{12} & \frac{1}{6} \\
\hline
8 & 1 & 0 & \frac{1}{3} & -\frac{1}{2}
\end{array}
\end{array}$$

## Simplex Method -

Q. Solve the following LPP

$$\text{Max } Z = 80x_1 + 55x_2$$

Subject to

$$4x_1 + 2x_2 \leq 40$$

$$2x_1 + 4x_2 \leq 32$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① SLPP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

$$4x_1 + 2x_2 + s_1 = 40$$

$$2x_1 + 4x_2 + s_2 = 32$$

and  $x_1, x_2, s_1, s_2 \geq 0$

Step② Represent in Matrix form

$$AX = B$$

$$\begin{bmatrix} x_1 & x_2 & s_1 & s_2 \\ 4 & 2 & 1 & 0 \\ 2 & 4 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 40 \\ 32 \end{bmatrix}$$

Step③ Construct starting simplex table

Basic Variable	C <sub>B</sub>	C <sub>j</sub>	80	55	0	0	MinnR.
	x <sub>B</sub>	x <sub>1</sub>	x <sub>2</sub>	s <sub>1</sub>	s <sub>2</sub>	x <sub>B</sub> /x <sub>i</sub>	
← S <sub>1</sub>	0	40	4	2	1	0	40/4=10
S <sub>2</sub>	0	32	2	4	0	1	32/2=16
Z = 0		A <sub>j</sub> → -80	-55	0	0		

outgoing vector

(x<sub>2</sub>) incoming vector.

Basic Variable	C <sub>B</sub>	C <sub>j</sub>	80	55	0	0	mink
	x <sub>B</sub>	x <sub>1</sub>	x <sub>2</sub>	s <sub>1</sub>	s <sub>2</sub>	x <sub>B</sub> /x <sub>i</sub>	
x <sub>1</sub>	80	10	1	1/2	1/4	0	20
← S <sub>2</sub>	0	12	6	3	-1/2	1	4 ←

outgoing vector

incoming

Basic Variable	C <sub>B</sub>	C <sub>j</sub>	80	55	0	0
	x <sub>B</sub>	x <sub>1</sub>	x <sub>2</sub>	s <sub>1</sub>	s <sub>2</sub>	
x <sub>1</sub>	80	8	1	0	1/3	-1/6
x <sub>2</sub>	55	4	0	1	-1/6	1/3

$$Z = 860$$

$$A_j \rightarrow 0 \quad 0 \quad 35/2 \quad 5$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: Step① S LPP

$$\text{Max } Z = 80x_1 + 55x_2 + 0s_1 + 0s_2$$

Subject to

$$4x_1 + 2x_2 + s_1 = 40$$

$$2x_1 + 4x_2 + s_2 = 32$$

outgoing vector

Basic Variable	C <sub>B</sub>	x <sub>1</sub>	x <sub>2</sub>	s <sub>1</sub>	s <sub>2</sub>	Z <sub>0</sub>
x <sub>1</sub>	80	1	0	0	0	40

Simplex Method | Minimization problem | operational research #10

## Minimization

## Simplex Method

Solution of LPP



By:-  
Sandeep Kumar Gour

Basic Va	x <sub>1</sub>	x <sub>2</sub>	s <sub>1</sub>	s <sub>2</sub>	Z <sub>0</sub>
x <sub>1</sub>	80	0	1	0	40
x <sub>2</sub>	55	4	0	1	32

Aj	0	0	35/2	5
Z = 860				

General LPP to Standard LPP with examples in hindi  
(lecture 08)

## LPP

## General $\rightarrow$ Standard

By:-  
Sandeep Kumar Gour



since all  $\Delta_j \geq 0$ , optimal basic feasible sol<sup>n</sup> is obtained.

$$\text{Max } Z = 860, x_1 = 8, x_2 = 4.$$

## Simplex Method - Minimization Problem

Ques: GLPP : Min  $Z = x_1 - 3x_2 + 2x_3$   
subject to

$$3x_1 - x_2 + 3x_3 \leq 7$$

$$-2x_1 + 4x_2 \leq 12$$

$$-4x_1 + 3x_2 + 8x_3 \leq 10$$

and  $x_1, x_2, x_3 \geq 0$

Sol<sup>n</sup>: SLPP: Max  $Z' = -x_1 + 3x_2 - 2x_3$   
subject to

$$3x_1 - x_2 + 3x_3 + s_1 = 7$$

$$-2x_1 + 4x_2 + s_2 = 12$$

$$-4x_1 + 3x_2 + 8x_3 + s_3 = 10$$

and  $x_1, x_2, x_3, s_1, s_2, s_3 \geq 0$

## Simplex Method - Minimization Problem

Ques: GLPP : Min  $Z = x_1 - 3x_2 + 2x_3$   
subject to

$$3x_1 - x_2 + 3x_3 \leq 7$$

$$-2x_1 + 4x_2 \leq 12$$

$$-4x_1 + 3x_2 + 8x_3 \leq 10$$

and  $x_1, x_2, x_3 \geq 0$

Sol<sup>n</sup>: SLPP: Max  $Z' = -x_1 + 3x_2 - 2x_3 + 0s_1 + 0s_2 + 0s_3$   
subject to

$$3x_1 - x_2 + 3x_3 + s_1 = 7$$

$$-2x_1 + 4x_2 + s_2 = 12$$

$$-4x_1 + 3x_2 + 8x_3 + s_3 = 10$$

and  $x_1, x_2, x_3, s_1, s_2, s_3 \geq 0$

SOL. • SLPP: Max  $Z^I = -x_1 + 3x_2 - 2x_3 + 0s_1 + 0s_2 + 0s_3$

Subject to

$$3x_1 - x_2 + 3x_3 + s_1 = 7$$

$$-2x_1 + 4x_2 + s_2 = 12$$

$$-4x_1 + 3x_2 + 8x_3 + s_3 = 10$$

and  $x_1, x_2, x_3, s_1, s_2, s_3 \geq 0$

Representation in Matrix form

$$AX = B$$

$$\begin{bmatrix} x_1 & x_2 & x_3 & s_1 & s_2 & s_3 \\ 3 & -1 & 3 & 1 & 0 & 0 \\ -2 & 4 & 0 & 0 & 1 & 0 \\ -4 & 3 & 8 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ s_1 \\ s_2 \\ s_3 \end{bmatrix} = \begin{bmatrix} 7 \\ 12 \\ 10 \end{bmatrix}$$

$$x_1 = x_2 = x_3 = 0$$

$$s_1 = 7, \quad s_2 = 12, \quad s_3 = 10$$

Initial Basic feasible soln

construct starting simplex table

	$c_j$	-1	3	-2	0	0	0	
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
$s_1$	0	7	3	-1	3	1	0	0
$s_2$	0	12	-2	4	0	0	1	0
$s_3$	0	10	-4	3	8	0	0	1
$Z=0$		$A_j^T$	1	-3	2	0	0	0

construct starting simplex table

	$C_j$	-1	3	-2	0	0	0		
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$	Min Ratio
$s_1$	0	7	3	-1	3	1	0	0	$7/1 = 7 \leftarrow$
$s_2$	0	12	-2	4	0	0	1	0	$12/4 = 3 \leftarrow$
$s_3$	0	10	-4	3	8	0	0	1	$10/3 = 3.3..$
$Z=0$			$\Delta_j$	1	-3↑	2	0	0	

↑ incoming vector

construct starting simplex table

	$c_j$	-1	3	-2	0	0	0	
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
$s_1$	0	7	3	-1	3	1	0	0
$s_2$	0	12	-2	4	0	0	1	0
$s_3$	0	10	-4	3	8	0	0	1

$Z = 0$        $\Delta_j$       1      -3↑      2      0      0      0

Min Ratio  
 $x_B/x_k, x_k > 0$

←  
 outgoing vector  
 if  
 $s_2$   
 $s_3$

$12/4 = 3 \leftarrow$   
 $10/3 = 3.3\dots$

↑ incoming vector

	$c_j$							
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$

$S_1$	0	7	3	-1	3	1	0	0
$S_2$	0	(12)	-2	4	0	0	1	0
$S_3$	0	10	-4	3	8	0	0	1
	$Z=0$	$A_j$	1	-3↑	2	0	0	0

↑ outgoing vector      ↓ incoming vector

$12/4 = 3 \leftarrow$        $10/3 = 3.3\dots$

		$C_j$	-1	3	-2	0	0	0
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$x_3$	$S_1$	$S_2$	$S_3$
$S_1$	0							
$x_2$	3	3		$-\frac{1}{2}$	1	0	0	$y_4$
$S_3$	0							

$$R_1 = R_1 + -\frac{1}{2}R_2$$

$$\begin{array}{ccccccc|ccc}
& & & 7 & 3 & -1 & 3 & 1 & 0 & 0 \\
& & & -\frac{1}{2} & 1 & 0 & 0 & 0 & & \\
\hline
& & & 3 & 1 & & & & & 
\end{array}$$

output vector

$S_1$	0	7	3	-1	3	1	0	0
$S_2$	0	(12)	-2	4	0	0	1	0
$S_3$	0	10	-4	3	8	0	0	1
	$Z=0$	$A_j$	1	-3↑	2	0	0	0

12/4 = 3 ←

10/3 = 3.3..

incoming vector

	$C_B$	$x_B$	$c_j$	-1	3	-2	0	0	0
Basic variable				$x_1$	$x_2$	$x_3$	$S_1$	$S_2$	$S_3$
$S_1$	0	10		5/2	0	3	1	1/4	0
$x_2$	3	3		-1/2	1	0	0	1/4	0
$S_3$	0								

$$R_1 = R_1 + R_2$$

$$\begin{array}{r} 7 \quad 3 \quad -1 \quad 3 \quad 1 \quad 0 \quad 0 \\ 3 \quad -1/2 \quad 1 \quad 0 \quad 0 \quad 1/4 \quad 0 \\ \hline 10 \quad 5/2 \quad 0 \quad 3 \quad 1 \quad 1/4 \quad 0 \end{array}$$

	$S_1$	0	7	3	-1	3	1	0	0
	$S_2$	0	12	-2	4	0	0	1	0
	$S_3$	0	10	-4	3	8	0	0	1
$Z=0$			$\Delta_j$	1	-3↑	2	0	0	0

↑ incoming vector

$12/4 = 3 \leftarrow$

$10/3 = 3.3..$

	$c_j$	-1	3	-2	0	0	0	0
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
$S_1$	0	10	$5/2$	0	3	1	$1/4$	0
$x_2$	3	3	$-1/2$	1	0	0	$1/4$	0
$S_3$	0	1	$-5/2$	0	8	0	$-3/4$	1

$Z=9$

$\Delta_j$

	$S_1$	0	7	3	-1	3	1	0	0
	$S_2$	0	12	-2	4	0	0	1	0
	$S_3$	0	10	-4	3	8	0	0	1
		$Z = 0$	$\Delta_j$	1	-3↑	2	0	0	0

outgoing vector

$12/4 = 3 \leftarrow$

$10/3 = 3.3..$

	$c_j$	-1	3	-2	0	0	0		
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$x_3$	$S_1$	$S_2$	$S_3$	
$S_1$	0	10	5/2	0	3	1	1/4	0	
$x_2$	3	3	-1/2	1	0	0	1/4	0	
$S_3$	0	1	-5/2	0	B	0	-3/4	1	
		$Z = 9$	$\Delta_j$	-1/2↑	0	2	0	$3/4$	0

outgoing vector

MinRatio:  
 $x_8/X_k, X_k > 0$   
 $10/5/2 = 4 \leftarrow$

$10/5/2 = 4 \leftarrow$

Basic variable	$C_B$	$X_B$	$C_j$	-1	3	-2	0	0	0
$s_1$	0	10	$\frac{5}{2}$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
$x_2$	3	3	$-\frac{1}{2}$		1	0	0	$\frac{1}{4}$	0
$s_3$	0	1	$-\frac{5}{2}$		0	8	0	$-\frac{3}{4}$	1

$Z' = 9$        $\Delta_j = -\frac{1}{2} \uparrow$       0      2      0       $\frac{3}{4}$       0

outgoing vector      incoming vector

Min Ratio  
 $x_B/X_k, X_k > 0$   
 $10/s_2 = 4 \leftarrow$

Basic variable	$C_B$	$X_B$	$C_j$	-1	3	-2	0	0	0
$x_1$	-1	4		1	0	$\frac{6}{5}$	$\frac{2}{5}$	$\frac{1}{10}$	0
$x_2$	3	5		0	1	$\frac{3}{5}$	$\frac{1}{5}$	$\frac{3}{10}$	0
$s_3$	0	11		0	0	11	1	$-\frac{1}{2}$	1

$$R_3 = R_3 + \frac{5}{2}R_1$$

1	$-\frac{5}{2}$	0	8	0	$-\frac{3}{4}$	1
10	$\frac{5}{2}$	0	3	1	$\frac{1}{4}$	0
11	0	0	11	1	$-\frac{1}{2}$	1

vector

$x_2$	3	3	$-\frac{1}{2}$	1	0	0	$\frac{1}{4}$	0
$s_3$	0	1	$-\frac{5}{2}$	0	8	0	$-\frac{3}{4}$	1
$Z' = 9$	$\Delta_j$	$-\frac{1}{2} \uparrow$	0	2	0	$\frac{3}{4}$	0	

incoming vector

	$c_j$	-1	3	-2	0	0	0	
Basic Variable	$c_B$	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
$x_1$	-1	4	1	0	$\frac{6}{5}$	$\frac{2}{5}$	$\frac{1}{10}$	0
$x_2$	3	5	0	1	$\frac{3}{5}$	$\frac{1}{5}$	$\frac{3}{10}$	0
$s_3$	0	11	0	0	11	1	$-\frac{1}{2}$	1
$Z' = 11$	$\Delta_j$	0	0	$\frac{13}{5}$	$\frac{1}{5}$	$-\frac{4}{5}$	0	

↑ vector

$$\left[ \begin{array}{c|ccccc|c} & x_2 & 3 & 3 & -\frac{1}{2} & 1 & 0 & 0 & y_4 & 0 \\ & s_3 & 0 & 1 & -\frac{5}{2} & 0 & B & 0 & -\frac{3}{4} & 1 \end{array} \right]$$

$Z' = 9 \quad \Delta_j \quad -\frac{1}{2} \uparrow \quad 0 \quad 2 \quad 0 \quad \frac{3}{4} \quad 0$

incoming  
vector

	$C_j$	-1	3	-2	0	0	0	
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
$x_1$	-1	4	1	0	$\frac{6}{5}$	$\frac{2}{5}$	$\frac{1}{10}$	0
$x_2$	3	5	0	1	$\frac{3}{5}$	$\frac{1}{5}$	$\frac{3}{10}$	0
$s_3$	0	11	0	0	1	1	$-\frac{1}{2}$	1

$Z' = 11 \quad \Delta_j \quad 0 \quad 0 \quad \frac{13}{5} \quad \frac{1}{5} \quad -\frac{4}{5} \quad 0$

$\therefore$  all  $\Delta_j \geq 0$  optimal sol<sup>n</sup> is obtained.

optimal sol<sup>n</sup> is  $\min Z = -11, x_1 = 4, x_2 = 5$ .

row vector

$x_2$	3	3	$-\frac{1}{2}$	1	0	0	$\frac{1}{4}$	0
$s_3$	0	1	$-\frac{5}{2}$	0	8	0	$-\frac{3}{4}$	1
$Z' = 9$	$\Delta_j$	$-\frac{1}{2} \uparrow$	0	2	0	$\frac{3}{4}$	0	

incoming vector

		$c_j$	-1	3	-2	0	0	0
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$x_3$	$s_1$	$s_2$	$s_3$
$x_1$	-1	4	1	0	$\frac{6}{5}$	$\frac{2}{5}$	$\frac{1}{10}$	0
$x_2$	3	5	0	1	$\frac{3}{5}$	$\frac{1}{5}$	$\frac{3}{10}$	0
$s_3$	0	11	0	0	11	1	$-\frac{1}{2}$	1

$$Z' = 11 \quad \Delta_j \quad 0 \quad 0 \quad \frac{13}{5} \quad \frac{1}{5} \quad -\frac{4}{5} \quad 0$$

$\therefore$  all  $\Delta_j \geq 0$  optimal sol<sup>n</sup> is obtained.

optimal sol<sup>n</sup> is  $\min Z = 11, x_1 = 4, x_2 = 5, x_3 = 0$ .

## Simplex Method - Unbounded Solution

$$\text{Max } Z = 3x_1 + 4x_2$$

Subject to

$$x_1 - x_2 \leq 1$$

$$-x_1 + x_2 \leq 2$$

$$\text{and } x_1, x_2 \geq 0$$

Sol<sup>n</sup>. SLPP

$$\text{Max } Z = 3x_1 + 4x_2 + 0S_1 + 0S_2$$

Subject to

$$x_1 - x_2 + S_1 = 1$$

$$-x_1 + x_2 + S_2 = 2$$

$$\text{and } x_1, x_2, S_1, S_2 \geq 0$$

Initial Basic Feasible solution

$$x_1 = x_2 = 0 \quad [S_1 = 1, S_2 = 2]$$

Matrix form  $A\bar{x} = B$

$$\begin{bmatrix} 1 & -1 & 1 & 0 \\ -1 & 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ S_1 \\ S_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$$

## Simplex Method - Unbounded Solution

$$\text{Max } Z = 3x_1 + 4x_2$$

Subject to

$$x_1 - x_2 \leq 1$$

$$-x_1 + x_2 \leq 2$$

$$\text{and } x_1, x_2 \geq 0$$

Sol<sup>n</sup>: SLP

$$\text{Max } Z = 3x_1 + 4x_2 + 0s_1 + 0s_2$$

Subject to

$$x_1 - x_2 + s_1 = 1$$

$$-x_1 + x_2 + s_2 = 2$$

$$\text{and } x_1, x_2, s_1, s_2 \geq 0$$

Initial Basic Feasible solution

$$x_1 = x_2 = 0 \quad \boxed{s_1 = 1, s_2 = 2}$$

Matrix form

$$\begin{bmatrix} 1 & -1 & 1 \\ -1 & 1 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$$

construct starting simplex table.

	$c_j$	3	4	0	0	
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$
$s_1$	0	1	1	-1	1	0
$s_2$	0	2	-1	1	0	1
Z=0		$\alpha_j$	-3	-4	0	0

## Simplex Method- Unbounded Solution

$$\text{Max } Z = 3x_1 + 4x_2$$

Subject to

$$x_1 - x_2 \leq 1$$

$$-x_1 + x_2 \leq 2$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup>: SLPP

$$\text{Max } z = 3x_1 + 4x_2 + 0s_1 + 0s_2$$

Subject to

$$x_1 - x_2 + s_1 = 1$$

$$-x_1 + x_3 + 5_2 = 2$$

and  $x_1, x_2, s_1, s_2 \geq 0$

#### Initial Basic Feasible sol.

$$x_1 = x_2 = 0 \quad | \quad s_1 = 1, s_2 =$$

## Matrix Form

卷之二

construct starting simplex table

	$C_j$	3	4	0	0		Min Ratio
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$X_B/X_K, X_K > 0$
$s_1$	0	1	1	-1	1	0	-
$\leftarrow s_2$	0	(2)	-1	1	0	1	$2/1 \leftarrow$
$Z = 0$		$\Delta_j$	-3	-4↑	0	0	

$(x_k)$  incoming vector

## Simplex Method - Unbounded Solution

$$\text{Max } Z = 3x_1 + 4x_2$$

Subject to

$$x_1 - x_2 \leq 1$$

$$-x_1 + x_2 \leq 2$$

$$\text{and } x_1, x_2 \geq 0$$

Sol<sup>n</sup>: SLPP

$$\text{Max } Z = 3x_1 + 4x_2 + 0s_1 + 0s_2$$

Subject to

$$x_1 - x_2 + s_1 = 1$$

$$-x_1 + x_2 + s_2 = 2$$

$$\text{and } x_1, x_2, s_1, s_2 \geq 0$$

Initial Basic Feasible solution

$$x_1 = x_2 = 0 \quad [s_1 = 1, s_2]$$

Matrix form  $Ax =$

$$\begin{bmatrix} 1 & -1 & 1 & 0 \\ -1 & 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix}$$

constructed starting simplex table

	$c_j$	3	4	0	0	
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$
$s_1$	0	1	1	-1	1	0
$s_2$	0	2	-1	1	0	1

$Z = 0$

$\Delta_j = -3 \quad -4 \uparrow \quad 0 \quad 0$

(Incoming vector)

outgoing vector

Min Ratio  
 $x_B/x_k, x_k/x_j$

	$c_j$	3	4	0	0	
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$
$s_1$	0	3	0	0	1	1
$x_2$	4	2	-1	1	0	1

$$R_1 = R_1 + R_2 (\text{New})$$

$$\begin{array}{r} 1 \quad 1 \quad -1 \quad 1 \quad 0 \\ + 2 \quad -1 \quad 1 \quad 0 \quad 1 \\ \hline R_1 (\text{New}) \quad 3 \quad 0 \quad 0 \quad 1 \quad 1 \end{array}$$

## Simplex Method - Unbounded Solution

$$\text{Max } Z = 3x_1 + 4x_2$$

Subject to

$$x_1 - x_2 \leq 1$$

$$-x_1 + x_2 \leq 2$$

$$\text{and } x_1, x_2 \geq 0$$

Sol<sup>n</sup>: SLPP

$$\text{Max } Z = 3x_1 + 4x_2 + 0s_1 + 0s_2$$

Subject to

$$x_1 - x_2 + s_1 = 1$$

$$-x_1 + x_2 + s_2 = 2$$

$$\text{and } x_1, x_2, s_1, s_2 \geq 0$$

Initial Basic Feasible solution

$$x_1 = x_2 = 0 \quad [s_1 = 1, s_2 = 2]$$

Matrix form  $Ax = B$

$$\begin{bmatrix} 1 & -1 & 1 & 0 \\ -1 & 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$$

construct starting simplex table

	$c_j$	3	4	0	0	
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$
$s_1$	0	1	1	-1	1	0
$s_2$	0	2	-1	1	0	1

$Z = 0$        $\Delta_j = -3 - 4 \uparrow$       0 0  
 outgoing vector      incoming vector      2/1  $\leftarrow$

	$c_j$	3	4	0	0	
Basic variable	$c_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$
$s_1$	0	3	0	0	1	1
$x_2$	4	2	-1	1	0	

$Z = 8$        $\Delta_j = -\frac{7}{2} \uparrow$       0 0 4

## Simplex Method - Unbounded Solution

$$\text{Max } Z = 3x_1 + 4x_2$$

Subject to

$$x_1 - x_2 \leq 1$$

$$-x_1 + x_2 \leq 2$$

$$\text{and } x_1, x_2 \geq 0$$

Sol<sup>n</sup>: SLPP

$$\text{Max } Z = 3x_1 + 4x_2 + 0s_1 + 0s_2$$

Subject to

$$x_1 - x_2 + s_1 = 1$$

$$-x_1 + x_2 + s_2 = 2$$

$$\text{and } x_1, x_2, s_1, s_2 \geq 0$$

Initial Basic Feasible solution

$$x_1 = x_2 = 0 \quad [s_1 = 1, s_2 = 2]$$

Matrix form  $Ax = B$

$$\begin{bmatrix} 1 & -1 & 1 & 0 \\ -1 & 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$$

construct starting simplex table

	$C_j$	3	4	0	0		Min Ratio
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$X_B/X_k, X_k$
$s_1$	0	1	1	-1	1	0	-
$s_2$	0	2	-1	1	0	1	2/1 ←
		$Z=0$					
		$\Delta_j$	-3	-4↑	0	0	

outgoing vector

$Z=0$

$\Delta_j$

$(x_2)$  incoming vector

	$C_j$	3	4	0	0		
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	
$s_1$	0	3	0	0	1	1	
$x_2$	4	2	-1	1	0	1	
		$Z=8$					
		$\Delta_j$	-7↑	0	0	4	

: all elements of the column  $x_1$  are negative or zero then the solution is unbounded.

## Degeneracy in Simplex method ( Tie for minimum Ratio ) -

Q. Given LPP:  $\text{Max } Z = 2x_1 + x_2$

subject to  $4x_1 + 3x_2 \leq 12$

$$4x_1 + x_2 \leq 8$$

$$4x_1 - x_2 \leq 8$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup> SLPP:  $\text{Max } Z = 2x_1 + x_2 + 0s_1 + 0s_2 + 0s_3$

subject to

$$4x_1 + 3x_2 + s_1 = 12$$

$$4x_1 + x_2 + s_2 = 8$$

$$4x_1 - x_2 + s_3 = 8$$

and  $x_1, x_2, s_1, s_2, s_3 \geq 0$

Initial Basic Feasible solution (IBFS)

$$x_1 = x_2 = 0, s_1 = 12, s_2 = 8, s_3 = 8, Z = 0$$

Subject to  $4x_1 + 3x_2 \leq 12$

$$4x_1 + x_2 \leq 8$$

$$4x_1 - x_2 \leq 8$$

and  $x_1, x_2 \geq 0$

Sol<sup>n</sup> SLPP: Max  $Z = 2x_1 + x_2 + 0s_1 + 0s_2 + 0s_3$

Subject to

$$4x_1 + 3x_2 + s_1 = 12$$

$$4x_1 + x_2 + s_2 = 8$$

$$4x_1 - x_2 + s_3 = 8$$

and  $x_1, x_2, s_1, s_2, s_3 \geq 0$

Initial Basic Feasible solution (IBFS)

$$x_1 = x_2 = 0, s_1 = 12, s_2 = 8, s_3 = 8, Z = 0$$

Matrix form  $AX = B$

$$\begin{bmatrix} 4 & 3 & 1 & 0 & 0 \\ 4 & 1 & 0 & 1 & 0 \\ 4 & -1 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ s_1 \\ s_2 \\ s_3 \end{bmatrix} = \begin{bmatrix} 12 \\ 8 \\ 8 \end{bmatrix}$$

	$C_j$	2	1	0	0	0		Min Ratio		
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	$X_B/X_K$	$s_1/X_K$	$s_2/X_K$
$s_1$	0	12	4	3	1	0	0	$12/4=3$	—	—
$s_2$	0	8	4	1	0	1	0	$8/4=2$	$0/4=0$	$1/4=\frac{1}{4}$
$\leftarrow s_3$	0	8	4	-1	0	0	1	$8/4=2$	$0/4=0$	$0/4=0$
$Z=0$		$\Delta_j$	-2	-1	0	0	0			
$s_1$	0	4	0	4	1	0	-1	$4/4=1$	$R_1=R_1-4R_3$ (New)	
$\leftarrow s_2$	0	0	0	12	0	1	-1	$0/2=0$	$R_2=R_2-4R_3$ (New)	
$x_1$	2	2	1	-1/4	0	0	1/4	—		
$Z=4$		$\Delta_j$	0	-3/2	0	0	1/2			
$\leftarrow s_1$	0	4	0	0	1	-2	1/1	$4/1=\frac{4}{1}$	$R_1=R_1-4R_2$ (New)	
$x_2$	1	0	0	1	0	1/2	-1/2			
$x_1$	2	2	1	0	0	1/8	1/8	$2/1/8=16$	$R_3=R_3+\frac{1}{4}R_2$	
$Z=4$		$\Delta_j$	0	0	0	3/4	-1/4			
$s_3$	0	4	0	0	1	-2	1			
$x_2$	1	2	0	1	1/2	-1/2	0		$R_2=R_2+\frac{1}{2}R_1$	
$x_1$	2	3/2	1	0	-1/8	3/8	0		$R_3=R_3-\frac{1}{8}R_1$	
$Z=5$		$\Delta_j$	0	0	1/4	1/4	0			

$s_1$	0	12	4	3	1	0	0	$\frac{12}{4} = 3$	-	-
$s_2$	0	8	4	1	0	1	0	$\frac{8}{4} = 2$	$0/4 = 0$	$\frac{1}{4} = \frac{1}{4}$
$\leftarrow s_3$	0	8	4	-1	0	0	1	$\frac{8}{4} = 2$	$0/4 = 0$	$0/4 = 0$
$Z=0$	$\Delta_j$	$\frac{-2}{(x_1)}$	-1	0	0	0	0			
$s_1$	0	4	0	4	1	0	-1	$\frac{4}{4} = 1$	$R_1 = R_1 - 4R_3$ (New)	
$\leftarrow s_2$	0	0	0	12	0	1	-1	$0/2 = 0$	$R_2 = R_2 - 4R_3$ (New)	
$x_1$	2	2	1	$-\frac{1}{4}$	0	0	$\frac{1}{4}$	-		
$Z=4$	$\Delta_j$	0	$-\frac{3}{2}$	1	0	0	$\frac{1}{2}$			
$\leftarrow s_1$	0	4	0	0	1	-2	1	$\frac{4}{1} = 4$	$R_1 = R_1 - 4R_2$ (New)	
$x_2$	1	0	0	1	0	$\frac{1}{2}$	$-\frac{1}{2}$			
$x_1$	2	2	1	0	0	$\frac{1}{8}$	$\frac{1}{8}$	$2/\frac{1}{8} = 16$	$R_3 = R_3 + \frac{1}{4}R_2$	
$Z=4$	$\Delta_j$	0	0	0	$\frac{3}{4}$	$-\frac{1}{4}$	1			
$s_3$	0	4	0	0	1	-2	1			$R_2 = R_2 + \frac{1}{2}R_1$
$x_2$	1	2	0	1	$\frac{1}{2}$	$-\frac{1}{2}$	0			$R_3 = R_3 - \frac{1}{8}R_1$
$x_1$	2	$\frac{3}{2}$	1	0	$-\frac{1}{8}$	$\frac{3}{8}$	0			
$Z=5$	$\Delta_j$	0	0	$\frac{1}{4}$	$\frac{1}{4}$	0				

$\therefore$  all  $\Delta_j \geq 0$ , optimal soln is obtained.

Optimal soln Max  $Z=5$ ,  $x_1 = \frac{3}{2}$ ,  $x_2 = 2$ .

## Two-phase Method-

Given GLPP  $\text{Max } Z = 2x_1 - x_2$

$$\text{Max } Z = 2x_1 + x_2$$

s.t.

$$4x_1 + 3x_2 \leq 12 \Rightarrow 4\underline{x_1} + 3\underline{x_2} + s_1 = 12$$

$$4x_1 + 2x_2 \leq 8 \Rightarrow 4x_1 + 2\underline{x_2} + s_2 = 8$$

$$\text{and } x_1, x_2 \geq 0 \quad \underline{x_1 = x_2 = 0} \quad \boxed{s_1 = 12, s_2 = 8}$$

$$\text{Max } Z = -2x_1 - x_2$$

s.t.

$$3x_1 + x_2 = 3 \Rightarrow 3x_1 + x_2 \stackrel{x}{=} 3$$

$$4x_1 + 3x_2 \geq 6 \Rightarrow 4x_1 + 3x_2 - s_1 = 6 \Rightarrow s_1 = -6$$

$$x_1 + 2x_2 \leq 6 \Rightarrow x_1 + 2x_2 + s_2 = 6$$

$$\text{and } x_1, x_2 \geq 0$$

G Max Z = -2x<sub>1</sub> - x<sub>2</sub>

s.t.  $3x_1 + x_2 = 3 \Rightarrow 3x_1 + x_2 = 3$

 $4x_1 + 3x_2 \geq 6 \Rightarrow 4x_1 + 3x_2 - s_1 = 6$ 
 $x_1 + 2x_2 \leq 6 \Rightarrow x_1 + 2x_2 + s_2 = 6$

and x<sub>1</sub>, x<sub>2</sub> ≥ 0

artificial

≤ + slack

≥ - surplus + artificial

= + artificial

I: eliminate

II: optimal soln

## Two-phase Method-

Given GLPP     $\text{Max } Z = 3x_1 - x_2$

Subject to the constraints

$$2x_1 + x_2 \geq 2 \quad \dots \quad ①$$

$$x_1 + 3x_2 \leq 2 \quad \dots \quad ②$$

$$x_2 \leq 4 \quad \dots \quad ③$$

and  $x_1, x_2 \geq 0$

s    sol<sup>n</sup>    SLPP     $\text{Max } Z = 3x_1 - x_2$

s.t.

$$2x_1 + x_2 - s_1 + a_1 = 2$$

$$x_1 + 3x_2 + s_2 = 2$$

$$x_2 + s_3 = 4$$

and  $x_1, x_2, s_1, s_2, s_3, a_1 \geq 0$

## Phase-I: Eliminating artificial variables

The simplex method is applied to a specially constructed auxiliary linear programming problem (ALPP).

Step 1: Assign a cost -1 to each artificial variable and a cost 0 to all other variables in the objective function.

Step 2: Construct the Auxiliary LPP in which the new objective function  $Z^*$  is to be maximized subject to the constraints.

Step 3: Solve the auxiliary problem by simplex method until either of the following three possibilities do arise.

- (i)  $\text{Max } Z^* < 0$  and atleast one artificial vector appear in the optimum basis at a positive level, then given problem does not possess any feasible solution.

S      sol<sup>n</sup>    SLPP

. Max  $Z = 3x_1 - x_2$

s.t.

$$2x_1 + x_2 - s_1 + a_1 = 2$$

$$x_1 + 3x_2 + s_2 = 2$$

$$x_2 + s_3 = 4$$

and  $x_1, x_2, s_1, s_2, s_3, a_1 \geq 0$

Auxiliary linear programming problem

$$\text{Max } Z^* = 0x_1 + 0x_2 + 0s_1 + 0s_2 + 0s_3 - a_1$$

s.t.

$$2x_1 + x_2 - s_1 + a_1 = 2$$

$$x_1 + 3x_2 + s_2 = 2$$

$$x_2 + s_3 = 4$$

and  $x_1, x_2, s_1, s_2, s_3, a_1 \geq 0$

Phase-I

	$C_j$	0	0	0	0	0	-1		
Basic Variable	$C_B$	$X_B$	$x_1$	$x_2$	$S_1$	$S_2$	$S_3$	$a_1$	Min Ratio $X_B/X_R$
$a_1$	-1	2	2	1	-1	0	0	1	
$S_2$	0	2	1	3	0	1	0	0	
$S_3$	0	4	0	1	0	0	1	0	
$Z^* = -2$			$\Delta_j$	-2	-1				

$$Z^* = C_B X_B = (-1, 0, 0)(2, 2, 4)$$

$$= -2 + 0 + 0 = -2$$

$$\Delta_j = C_B X_j - C_j \Rightarrow \Delta_1 = (-1, 0, 0)(2, 1, 0) - 0$$

$$= -2 + 0 + 0 - 0 = -2$$

$$\Delta_2 = C_B X_2 - C_2 = (-1, 0, 0)(1, 3, 1) - 0$$

$$= -1 + 0 + 0 - 0 = -1$$

Phase-I

	$C_j$	0	0	0	0	0	-1	
Basic Variable	$C_B$	$x_B$	$x_1$	$x_2$	$S_1$	$S_2$	$S_3$	$a_1$
$\leftarrow a_1$	-1	2	2	1	-1	0	0	1
$S_2$	0	2	1	3	0	1	0	0
$S_3$	0	4	0	1	0	0	1	0
$Z^* = -2$		$\Delta j$	$\frac{-2}{(x_1)} \uparrow$	-1	1	0	0	0

Min Ratio  
 $x_B/x_K, x_W > 0$

$$2/2 = k$$

$$2/1 = 2$$

### phase-I

	$C_j$	0	0	0	0	0	-1	
Basic variable	$C_B$	$x_B$	$x_1$	$x_2$	$S_1$	$S_2$	$S_3$	$a_1$
$\leftarrow$ outgoing $a_1$	-1	2	2	1	-1	0	0	1
$S_2$	0	2	1	3	0	1	0	0
$S_3$	0	4	0	1	0	0	1	0
$Z^* = -2$		$\Delta_j$	$\frac{-2}{x_1} \uparrow -1$ $(x_k)$ incoming	1	0	0	0	
$x_1$	0	1	2	$1/2$	$-1/2$	0	0	$1/2$
$S_2$	0	1	0	$5/2$	$1/2$	1	0	$-1/2$
$S_3$	0	4	0	1	0	0	1	0
$Z^* = 0$		$\Delta_j$	0	0	0	0	0	1

Min Ratio  
 $x_B/x_K, x_K > 0$

$$2/2 = 1$$

$$2/1 = 2$$

Step 1: Assign a cost -1 to each artificial variable and a cost 0 to all other variables in the objective function.

Step 2: Construct the Auxiliary LPP in which the new objective function  $Z^*$  is to be maximized subject to the constraints.

Step 3: Solve the auxiliary problem by simplex method until either of the following three possibilities do arise.

- (i)  $\text{Max } Z^* < 0$  and atleast one artificial vector appear in the optimum basis at a positive level, then given problem does not possess any feasible solution.
- (ii)  $\text{Max } Z^* = 0$  and at least one artificial vector appears in the optimum basis at a zero level. then proceed to phase II.
- (iii)  $\text{Max } Z^* = 0$  and no one artificial vector appears in the optimum basis, then proceed to phase II.

### Phase-I

	$C_j$	0	0	0	0	0	-1	
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	$a_1$
$a_1$ outgoing	-1	2	2	1	-1	0	0	1
$s_2$	0	2	1	3	0	1	0	0
$s_3$	0	4	0	1	0	0	1	0
$Z^* = -2$		$\Delta_j$	$\frac{-2}{(x_1)} \uparrow -1$	1	0	0	0	
$x_1$	0	1	$\frac{1}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$	0	0	$\frac{1}{2}$
$s_2$	0	1	0	$\frac{5}{2}$	$\frac{1}{2}$	1	0	$-\frac{1}{2}$
$s_3$	0	4	0	1	0	0	1	0
$Z^* = 0$		$\Delta_j$	0	0	0	0	0	1

Min Ratio  
 $x_B/X_K, X_B > 0$

$$2/2 = 1 \leftarrow \\ 2/1 = 2$$

subject to the constraints.

Step 3: Solve the auxiliary problem by simplex method until either of the following three possibilities do arise.

- (i)  $\text{Max } Z^* < 0$  and atleast one artificial vector appear in the optimum basis at a positive level, then given problem does not possess any feasible solution.
- (ii)  $\text{Max } Z^* = 0$  and at least one artificial vector appears in the optimum basis at a zero level. then proceed to phase II.
- (iii)  $\text{Max } Z^* = 0$  and no one artificial vector appears in the optimum basis, then proceed to phase II.

Phase II: Get an optimal solution.

Now assign the actual cost to the variables in the objective function and a zero cost to every artificial variable that appears in the basis at the zero level.

			$Z$	(X)	Incoming				
$x_1$	0	1	21	$\frac{1}{2}$	$-\frac{1}{2}$	0	0	$\frac{1}{2}$	
$s_2$	0	1	0	$\frac{5}{2}$	$\frac{1}{2}$	1	0	$-\frac{1}{2}$	
$s_3$	0	4	0	1	0	0	1	0	
$Z = 0$			$\Delta_j$	0	0	0	0	0	1

$\therefore$  all  $\Delta_j \geq 0$ , Max  $Z = 0$ , No artificial vector appears in the basis, we proceed to phase II.

	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	$C_j$	MR
$x_1$	3	1	1	$\frac{1}{2}$	$-\frac{1}{2}$	0	0		$X_B/X_R$
$\leftarrow s_2$	0	1	0	$\frac{5}{2}$	$\frac{1}{2}$	1	0		
$\downarrow s_3$	0	4	0	1	0	0	1		
$Z = 3$			$\Delta_j$	0	$\frac{5}{2}$	$-\frac{3}{2}$	↑ 0	0	$1/\frac{1}{2} = 2 \leftarrow$

$$Z=0 \quad \Delta_j \mid 0 \ 0 \ 0 \ 0 \ 0 \ 1$$

$\because$  all  $\Delta_j \geq 0$ , Max  $Z=0$ , No artificial vector appears in the basis, we proceed to phase II.

Phase-II.

	$C_j$	3	-1	0	0	0		
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	M.R $X_B/X_k$
$x_1$	3	1	1	$\frac{1}{2}$	$\frac{1}{2}$	0	0	
$\leftarrow s_2$ outgoing	0	1	0	$\frac{s_2}{2}$	$\frac{1}{2}$	1	0	
$s_3$	0	4	0	1	0	0	1	
$Z=3$	$\Delta_j$	0	$\frac{s_2}{2}$	$-\frac{3}{2}$	0	0	0	
$x_1$	3	2	1	3	0	1	0	
$s_1$	0	2	0	5	1	2	0	
$s_3$	0	4	0	1	0	0	1	
$Z=6$	$\Delta_j$	0	10	0	3	0		

$\therefore$  all  $\Delta_j \geq 0$ , optimal sol<sup>n</sup> is obtained.

Optimal sol<sup>n</sup> Max  $Z=6$ ,  $x_1=2$ ,  $x_2=0$ .

## Infeasible (No feasible) Solution in Two-phase method-

$$\text{Min } Z = x_1 - 2x_2 - 3x_3$$

Subject to constraints

$$-2x_1 + x_2 + 3x_3 = 2$$

$$2x_1 + 3x_2 + 4x_3 = 1$$

and  $x_1, x_2, x_3 \geq 0$

Sol<sup>n</sup>: SLPP

$$\text{Max } Z^1 = -x_1 + 2x_2 + 3x_3$$

s.t.

$$-2x_1 + x_2 + 3x_3 + a_1 = 2$$

$$2x_1 + 3x_2 + 4x_3 + a_2 = 1$$

and  $x_1, x_2, x_3, a_1, a_2 \geq 0$

Auxiliary LPP

$$\text{Max } Z^* = 0x_1 + 0x_2 + 0x_3 - a_1 - a_2$$

$$\text{s.t. } -2x_1 + x_2 + 3x_3 + a_1 = 2$$

$$2x_1 + 3x_2 + 4x_3 + a_2 = 1$$

and  $x_1, x_2, x_3, a_1, a_2 \geq 0$

		$C_j$	0	0	0	-1	-1	
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$x_3$	$a_1$	$a_2$	
$a_1$	-1	2	-2	1	(3)	1	0	
$a_2$	-1	1	2	3	(4)	0	1	

$Z^* = -3$

$\Delta_j$

↑ (Y<sub>k</sub>) increasing

outgoing ←  $a_2$

Min Ratio  
 $x_B/x_k$   
2/3 ←

1/4 ←

## Infeasible (No feasible) Solution in Two-phase method-

$$\text{Min } Z = x_1 - 2x_2 - 3x_3$$

Subject to constraints

$$-2x_1 + x_2 + 3x_3 = 2$$

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and  $x_1, x_2, x_3 \geq 0$

Sol<sup>n</sup>: SLPP

$$\text{Max } Z^1 = -x_1 + 2x_2 + 3x_3$$

s.t.

$$-2x_1 + x_2 + 3x_3 + a_1 = 2$$

$$2x_1 + 3x_2 + 4x_3 + a_2 = 1$$

and  $x_1, x_2, x_3, a_1, a_2 \geq 0$

Auxiliary LPP

$$\text{Max } Z^* = 0x_1 + 0x_2 + 0x_3 - a_1 - a_2$$

$$\text{s.t. } -2x_1 + x_2 + 3x_3 + a_1 = 2$$

$$2x_1 + 3x_2 + 4x_3 + a_2 = 1$$

and  $x_1, x_2, x_3, a_1, a_2 \geq 0$

Basic variable	c_B	Cj						Min Ratio $x_B/x_{k_e}$
		0	0	0	-1	-1		
$a_1$	-1	2	-2	1	(3)	1	0	2/3
$a_2$	-1	1	2	3	(4)	0	1	1/4
$Z^* = -3$		$\Delta_j$	0	-4	-7	0	0	$(x_k)$ incoming
$a_1$	-1	5/4	-7/2	-5/4	0	1	-3/4	
$x_3$	0	1/4	1/2	3/4	1	0	1/4	
$Z^* = -5/4$		$\Delta_j$	7/2	5/4	0	0	7/4	

$\because$  all  $\Delta_j \geq 0$ , But  $\text{Max } Z^* = -5/4$  is negative and the artificial variable  $a_1$  appears in the basis solution at a positive level. Hence the original problem does not possess any feasible solution.

## Big-M method (Charnes' Penalty method)

### Steps:

1. Express the problem in the standard form.
2. Assign the large negative price -M to the artificial variables.
3. Use simplex method.

Qn: Given LPP

$$\text{Max } Z = -2x_1 - x_2$$

Subject to the constraints

$$3x_1 + x_2 = 3$$

$$4x_1 + 3x_2 \geq 6$$

$$x_1 + 2x_2 \leq 4$$

$$\text{and } x_1, x_2 \geq 0$$

## Big-M method (Charnes's Penalty method)

Steps:

1. Express the problem in the standard form.
2. Assign the large negative price  $-M$  to the artificial variables.
3. Use simplex method.

$$\begin{array}{ll} \text{Max} & -M \\ \text{Min} & +M \end{array}$$

Qn: Given LPP

$$\text{Max } Z = -2x_1 - x_2$$

Subject to the constraints

$$3x_1 + x_2 = 3$$

$$4x_1 + 3x_2 \geq 6$$

$$x_1 + 2x_2 \leq 4$$

$$\text{and } x_1, x_2 \geq 0$$

Q4: Given LPP

$$\text{Max } Z = -2x_1 - x_2$$

Subject to the constraints

$$3x_1 + x_2 = 3$$

$$4x_1 + 3x_2 \geq 6$$

$$x_1 + 2x_2 \leq 4$$

and  $x_1, x_2 \geq 0$

SLPP

$$\text{Max } Z = -2x_1 - x_2 + 0s_1 + 0s_2 - Ma_1 - Ma_2$$

s.t.

$$3x_1 + x_2 + a_1 = 3$$

$$4x_1 + 3x_2 - s_1 + a_2 = 6$$

$$x_1 + 2x_2 + s_2 = 4$$

and  $\underbrace{x_1, x_2}_{\geq 0}, \underbrace{s_1, s_2}_{\geq 0}, \underbrace{a_1, a_2}_{\geq 0}$

### Starting Simplex table

	$C_j^0$	-2	-1	0	0	-M	-M	Min Ratio	
Basic Variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$	$x_B/x_R$
$a_1$	-M	3	3	1	0	0	1	0	
$a_2$	-M	6	4	3	-1	0	0	1	
$s_2$	0	4	1	2	0	1	0	0	
$Z = -9M$			$\Delta_i$						

$$Z = C_B x_B = (-M, -M, 0)(3, 6, 4)$$

$$= -3M - 6M + 0 = -9M$$

$$\Delta_j^* = C_B x_j^* - C_j^*$$

$$\Delta_1 = C_B x_1 - C_1 = (-M, -M, 0)(3, 4, 1) - (-2)$$

$$= -3M - 4M + 0 + 2 = 2 - 7M$$

### Starting Simplex table

	$C_j^0$	-2	-1	0	0	-M	-M	Min Ratio	
Basic variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$	$x_B/x_K$
$a_1$	-M	3	3	1	0	0	1	0	
$a_2$	-M	6	4	3	-1	0	0	1	
$s_2$	0	4	1	2	0	1	0	0	
$Z = -9M$		$\Delta_j$	$2-7M$	$1-4M$	M	0	0	0	

$a_2$	-M	6	4	3	-1	0	0	0	7
$s_2$	0	4	1	2	0	1	0	0	4/1
$Z = -9M$	$\Delta_j^o$			$2-7M$	$1-4M$	$M$	0	0	

$\uparrow$  Incoming  
(Xk)

$$R_2 = R_2 - 4R_1 \quad R_3 = R_3 - R_1$$

	$C_B$	$C_j$	-2	-1	0	0	-M	-M	
Basic variable	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$		
$x_1$	-2	1	1	$\frac{1}{3}$	0	0	$\frac{1}{3}$	0	
$a_2$	-M	2	0	$\frac{5}{3}$	-1	0	$-\frac{4}{3}$	1	
$s_2$	0	3	0	$\frac{5}{3}$	0	1	$-\frac{1}{3}$	0	
	.	.	.	.	.	.	.	.	.

$$R_2 = R_2 - 4R_1 \quad R_3 = R_3 - R_1$$

	$C_j$	-2	-1	0	0	-M	-M	MR	
Basic variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$	$x_B/x_K$
$x_1$	-2	1	1	$\frac{1}{3}$	0	0	$\frac{1}{3}$	0	3
$\overset{\leftarrow}{a_2}$	$-M$	$\overset{\cancel{2}}{0}$	$\overset{\cancel{2}}{0}$	$\frac{5}{3}$	-1	0	$\frac{-4}{3}$	1	$\frac{6}{5} \leftarrow$
$s_2$	0	3	0	$\frac{5}{3}$	0	1	$-\frac{1}{3}$	0	$\frac{9}{5}$
$Z = -2 - 2M$	$A_j$	0	$\frac{1-5M}{3}$	M	0	$\frac{-2+7M}{3}$	0	.	.

$(x_K)$  incoming

	$C_j$	
Basic variable	$C_B$	$x_B$

$$\boxed{R_2 = R_2 - 4R_1} \quad \boxed{R_3 = R_3 - R_1}$$

	$C_j$	-2	-1	0	0	-M	-M	MR	
Basic variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$	$x_B/x_K$
$x_1$	-2	1	1	$\frac{1}{3}$	0	0	$\frac{1}{3}$	0	3
$\overset{\leftarrow \text{Outgoing}}{x_2}$	$\overset{\leftarrow \text{Outgoing}}{-M}$	$\overset{\leftarrow}{2}$	0	$\frac{5}{3}$	-1	0	$\frac{-4}{3}$	1	$\frac{6}{5} \leftarrow$
$s_2$	0	3	0	$\frac{5}{3}$	0	1	$-\frac{1}{3}$	0	$\frac{1}{5}$
$Z = -2 - M$	$\Delta_j$	0	$\frac{1-5M}{3}$	M	0	$\frac{-2+7M}{3}$	0		.
			( $x_K$ ) incoming						.

	$C_j$	-2	-1	0	0	-M	-M		
Basic variable	$C_B$	$x_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$	
$x_1$	-2	$\frac{3}{5}$	1	0	$\frac{1}{5}$	0	$\frac{3}{5}$	$-\frac{1}{5}$	$R_1 = R_1 - \frac{1}{3}R_2$
$x_2$	-1	$\frac{6}{5}$	0	1	$-\frac{3}{5}$	0	$-\frac{4}{5}$	$\frac{3}{5}$	
$s_2$	0	1	0	0	1	1	1	-1	$R_3 = R_3 - \frac{5}{3}R_2$
$Z = -\frac{12}{5}$	$\Delta_j$	0	0	$\frac{1}{5}$	0	$\frac{M-2}{5}$	$\frac{M-1}{5}$		.

variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$	$X_B/X_K$
$x_1$	-2	1	1	$\frac{1}{3}$	0	0	$\frac{1}{3}$	0	3
$\overset{\leftarrow}{\text{other}} a_2$	$-M$	$\frac{2}{3}$	0	$\frac{5}{3}$	-1	0	$-\frac{4}{3}$	1	$\frac{6}{5} \leftarrow$
$s_2$	0	3	0	$\frac{5}{3}$	0	1	$-\frac{1}{3}$	0	$\frac{1}{5}$
$Z = -2 - \frac{2M}{3}$	$\Delta_j$	0	$\frac{1-5M}{3}$	M	0	$\frac{-2+7M}{3}$	0	.	.

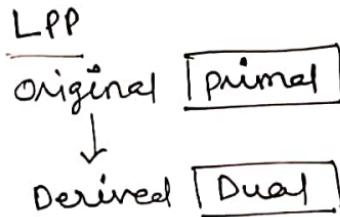
$(X_K)$  incoming

	$C_j$	-2	-1	0	0	$-M$	$-M$	
Basic variable	$C_B$	$X_B$	$x_1$	$x_2$	$s_1$	$s_2$	$a_1$	$a_2$
$x_1$	-2	$\frac{3}{5}$	1	0	$\frac{1}{5}$	0	$\frac{3}{5}$	$-\frac{1}{5}$
$x_2$	-1	$\frac{6}{5}$	0	1	$-\frac{3}{5}$	0	$-\frac{4}{5}$	$\frac{3}{5}$
$s_2$	0	1	0	0	1	1	1	-1
$Z = -\frac{12}{5}$	$\Delta_j$	0	0	$\frac{2}{5}$	0	$M - \frac{2}{5}$	$M - \frac{1}{5}$	.

$\therefore$  all  $\Delta_j \geq 0$ , Hence optimal sol<sup>n</sup> is obtained.

Optimal sol<sup>n</sup> Max  $Z = -\frac{12}{5}$ ,  $x_1 = \frac{3}{5}$ ,  $x_2 = \frac{6}{5}$ .

## Duality in Linear Programming



- ① [ Primal → Dual ]
- ② Dual Simplex method

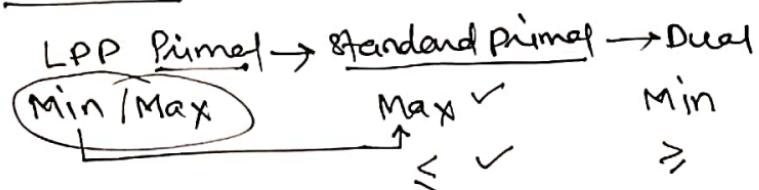
## Derived | Dual |

①  $\boxed{\text{Primal} \rightarrow \text{Dual}}$

1m.

② Dual simplex method

concept 1:



concept 2:

Primal  $\rightarrow$  Standard primal  $\rightarrow$  Dual

✓ Min  $\rightarrow$  Min  $\geq$  - Max  $\leq$   
Max  $\rightarrow$  Max  $\leq$                   Min  $\geq$

## Duality in Linear Programming

### Steps for a Standard primal Form :-

1. Change the objective function to Maximization form.
2. If the constraints have an 'inequality sign' ' $\geq$ '  
the convert into ' $\leq$ '
3. If the constraint has an '=' sign then replace  
it by two constraints involving the inequalities  
going in opposite directions.  
$$\text{Eg:- } x_1 + 2x_2 = 4 \Rightarrow x_1 + 2x_2 \leq 4 \checkmark$$
$$x_1 + 2x_2 \geq 4 \Rightarrow -x_1 - 2x_2 \leq -4$$

4. Every unrestricted variable is replaced by the  
difference of two non-negative variables.
5. Standard primal form of the given LPP in  
which -

i) All constraints have ' $\leq$ ' sign, where the  
objective function is of Maximization form.

## Duality in Linear Programming

### Steps for a Standard primal form :-

1. Change the objective function to Maximization form.
2. If the constraints have an 'inequality sign' ' $\geq$ '  
the convert into ' $\leq$ '
3. If the constraint has an '=' sign then replace  
it by two constraints involving the inequalities  
going in opposite directions.  
$$\text{Eg:- } x_1 + 2x_2 = 4 \Rightarrow x_1 + 2x_2 \leq 4 \checkmark$$
$$x_1 + 2x_2 \geq 4 \Rightarrow -x_1 - 2x_2 \leq -4$$

4. Every unrestricted variable is replaced by the  
difference of two non-negative variables.  $x_3 = x_3^1 - x_3^2$
5. Standard primal form of the given LPP in

which -

- i) All constraints have ' $\leq$ ' sign, where the  
objective function is of Maximization form.

1. Change the objective function to Maximization form.
2. If the constraints have an 'inequality sign' ' $\geq$ '  
the convert into ' $\leq$ '
3. If the constraint has an '=' sign then replace it by two constraints involving the inequalities going in opposite directions.

Eg:-  $x_1 + 2x_2 = 4 \Rightarrow x_1 + 2x_2 \leq 4$  ✓  
 $x_1 + 2x_2 \geq 4 \Rightarrow -x_1 - 2x_2 \leq -4$

4. Every unrestricted variable is replaced by the difference of two non-negative variables.  $x_3 = x_3^1 - x_3^2$
5. Standard primal form of the given LPP in which -

- ✓ (i) All constraints have ' $\leq$ ' sign, where the objective function is of Maximization form.
  - (ii) All constraints have ' $\geq$ ' sign, where the objective function is of Minimization form.
- Converting Primal into its Dual

## Rules for converting Primal into its Dual

1. Transpose the rows and columns of the constraint co-efficient.
2. Transpose the co-efficient of the objective function and the right side constants.
3. Change the inequalities from ' $\leq$ ' to ' $\geq$ ' sign.
4. Minimize the objective function instead of maximizing it.

### Standard primal

$$\text{Max } Z = 3x_1 + 4x_2$$

s.t.

$$2x_1 + 6x_2 \leq 16$$

$$-5x_1 - 2x_2 \leq -20$$

$$\text{and } x_1, x_2 \geq 0$$

### Dual

## Primal to Dual Conversion

$$\begin{array}{l} \textcircled{1} \text{ Max } Z = 3x_1 + 4x_2 \\ \text{s.t.} \\ 2x_1 + 6x_2 \leq 16 \quad \textcircled{1} \\ 5x_1 + 2x_2 \geq 20 \quad \textcircled{2} \end{array} \quad \left. \begin{array}{l} \text{primal} \\ \text{Max} \\ \leq \end{array} \right\}$$

and  $x_1, x_2 \geq 0$

### Standard primal

$$\text{Max } Z = [3x_1 + 4x_2]$$

s.t.

$$\begin{array}{l} 2x_1 + 6x_2 \leq 16 \\ -5x_1 - 2x_2 \leq -20 \end{array}$$

and  $x_1, x_2 \geq 0$

### Dual

$$\text{Min } Z = [16w_1 - 20w_2]$$

s.t.

$$\begin{array}{l} 2w_1 - 5w_2 \geq 3 \\ 6w_1 - 2w_2 \geq 4 \end{array}$$

and  $w_1, w_2 \geq 0$

$$\begin{array}{l} \boxed{x_1 + 6x_2 \leq 10} \\ -5x_1 - 2x_2 \leq -20 \\ \text{and } x_1, x_2 \geq 0 \end{array} \quad \left| \quad \begin{array}{l} \boxed{w_1 - 2w_2 \geq 4} \\ \text{and } w_1, w_2 \geq 0 \end{array} \right.$$

②  $\text{Min } Z = 2x_2 + 5x_3$  | standard primal  
 s.t.  
 $x_1 + x_2 \geq 2$   
 $2x_1 + x_2 + 6x_3 \leq 6$   
 $x_1 - x_2 + 3x_3 = 4$   
 and  $x_1, x_2, x_3 \geq 0$

$\text{Max } Z = -2x_2 - 5x_3$
s.t
$-x_1 - x_2 \leq -2$
$2x_1 + x_2 + 6x_3 \leq 6$
$x_1 - x_2 + 3x_3 \leq 4$
$-x_1 + x_2 - 3x_3 \leq -4$
and $x_1, x_2, x_3 \geq 0$

Dual  $\text{Min } Z = -2w_1 + 6w_2 + 4w_3 - 4w_4$

s.t.  $-w_1 + 2w_2 + w_3 - w_4 \geq 0$   
 $-w_1 + w_2 - w_3 + w_4 \geq -2$   
 $6w_2 + 3w_3 - 3w_4 \geq -5$   
 and  $w_1, w_2, w_3, w_4 \geq 0$

$$\textcircled{3} \quad \begin{aligned} \text{Min } Z &= x_1 + x_2 + x_3 \Rightarrow (-x_1 - x_2 - x_3) \\ \text{s.t. } x_1 - 3x_2 + 4x_3 &= 5 \rightarrow x_1 - 3x_2 + 4x_3 \leq 5 \\ x_1 - 2x_2 &\leq 3 \quad x_1 - 3x_2 + 4x_3 \geq 5 \\ 2x_2 - x_3 &\geq 4 \rightarrow 2x_2 - (x_3^1 - x_3^{11}) \geq 4 \end{aligned}$$

and  $x_1, x_2 \geq 0, x_3$  is unrestricted.  $|x_3 = x_3^1 - x_3^{11}|$

$$\stackrel{\text{Defn}}{=} \text{Max } Z = -x_1 - x_2 - (x_3^1 - x_3^{11})$$

$$\begin{aligned} \text{s.t. } x_1 - 3x_2 + 4(x_3^1 - x_3^{11}) &\leq 5 \\ -x_1 + 3x_2 - 4(x_3^1 - x_3^{11}) &\leq -5 \\ x_1 - 2x_2 &\leq 3 \\ -2x_2 + (x_3^1 - x_3^{11}) &\leq -4 \end{aligned}$$

$$-2x_2 + (x_3^1 - x_3^{11}) \leq -4$$

standard primal

$$\text{Max } Z = -x_1 - x_2 - x_3^1 + x_3^{11}$$

$$\text{s.t. } x_1 - 3x_2 + 4x_3^1 - 4x_3^{11} \leq 5$$

$$-x_1 + 3x_2 - 4x_3^1 + 4x_3^{11} \leq -5$$

$$x_1 - 2x_2 \leq 3$$

$$-2x_2 + x_3^1 - x_3^{11} \leq -4$$

$$\text{and } x_1, x_2, x_3^1, x_3^{11} \geq 0$$

$$-2x_2 + (x_3^1 - x_3^{11}) \leq -4$$

standard primal

$$\text{Max } Z = -x_1 - x_2 - x_3^1 + x_3^{11}$$

$$\text{s.t. } x_1 - 3x_2 + 4x_3^1 - 4x_3^{11} \leq 5$$

$$-x_1 + 3x_2 - 4x_3^1 + 4x_3^{11} \leq -5$$

$$x_1 - 2x_2 \leq 3$$

$$-2x_2 + x_3^1 - x_3^{11} \leq -4$$

$$\text{and } x_1, x_2, x_3^1, x_3^{11} \geq 0$$

$$x_1 - 2x_2 \leq 3$$

$$-2x_2 + (x_3' - x_3'') \leq -4$$

standard primal

$$\text{Max } Z = -x_4 - x_2 - x_3' + x_3''$$

$$\text{s.t. } x_1 - 3x_2 + 4x_3' - 4x_3'' \leq 5$$

$$-x_1 + 3x_2 - 4x_3' + 4x_3'' \leq -5$$

$$x_1 - 2x_2 + 0x_3' + 0x_3'' \leq 3$$

$$0x_1 - 2x_2 + x_3' - x_3'' \leq -4$$

$$\text{and } x_1, x_2, x_3', x_3'' \geq 0$$

Dual

$$\text{Min } Z = 5w_1 - 5w_2 + 3w_3 - 4w_4$$

$$\text{s.t. } w_1 - w_2 + w_3 \geq -1$$

$$-3w_1 + 3w_2 - 2w_3 - 2w_4 \geq -1$$

$$4w_1 - 4w_2 + w_4 \geq -1$$

$$-4w_1 + 4w_2 - w_4 \geq 1 \text{ and } w_1, w_2, w_3, w_4 \geq 0$$

## Primal to Dual Conversion -

### Concept-2

Primal	Standard	Dual
Min	Min $\geq$	Max $\leq$
Max	Max $\leq$	min $\geq$

find the dual?

Eg① Given Primal form

$$\text{Max } Z = 3x_1 + 10x_2 + 2x_3$$

Subject to

$$2x_1 + 3x_2 + 2x_3 \leq 7$$

$$3x_1 - 2x_2 + 4x_3 = 3$$

$$\text{and } x_1, x_2, x_3 \geq 0$$

$$3x_1 - 2x_2 + 4x_3 \leq 3 \leftarrow$$

$$(3x_1 - 2x_2 + 4x_3 \geq 3) \times -1$$

### standard primal

$$\text{Max } Z = 3x_1 + 10x_2 + 2x_3$$

subject to

$$2x_1 + 3x_2 + 2x_3 \leq 7$$

$$3x_1 - 2x_2 + 4x_3 = 3 \quad (3x_1 - 2x_2 + 4x_3 \geq 3) \times 1$$

and  $x_1, x_2, x_3 \geq 0$

standard primal

$$\text{Max } Z = 3x_1 + 10x_2 + 2x_3$$

subject to

$$2x_1 + 3x_2 + 2x_3 \leq 7$$

$$3x_1 - 2x_2 + 4x_3 \leq 3$$

$$-3x_1 + 2x_2 - 4x_3 \leq -3 \quad \text{and } x_1, x_2, x_3 \geq 0$$

Dual Form

$$\text{Min } Z = 7w_1 + 3w_2 - 3w_3$$

subject to

$$2w_1 + 3w_2 - 3w_3 \geq 3$$

$$3w_1 - 2w_2 + 2w_3 \geq 10$$

$$2w_1 + 4w_2 - 4w_3 \geq 2$$

$$\text{and } w_1, w_2, w_3 \geq 0$$



$$\begin{array}{l}
 \text{Min } Z = 3x_1 + 5x_2 + 6x_3 \\
 \text{Subject to} \\
 x_1 + 4x_2 + 6x_3 \leq 5 \\
 2x_1 + 3x_2 + 5x_3 \geq 4 \\
 3x_1 + x_2 + 7x_3 = 3 \\
 \text{and } x_1, x_2 \geq 0, \\
 x_3 \text{ is unrestricted.}
 \end{array}
 \quad
 \begin{array}{l}
 \text{Min } Z = 3x_1 + 5x_2 + 6(x_3^1 - x_3^{11}) \\
 \text{Subject to} \\
 -x_1 - 4x_2 - 6(x_3^1 - x_3^{11}) \geq -5 \\
 2x_1 + 3x_2 + 5(x_3^1 - x_3^{11}) \geq 4 \\
 3x_1 + x_2 + 7(x_3^1 - x_3^{11}) \geq 3 \\
 3x_1 + x_2 + 7(x_3^1 - x_3^{11}) \leq 3 \\
 -3x_1 - x_2 - 7(x_3^1 - x_3^{11}) \geq -3
 \end{array}$$

$$\text{Min } Z = 3x_1 + 5x_2 + 6x_3^1 - 6x_3^{11}$$

s.t.

$$\begin{aligned}
 -x_1 - 4x_2 - 6x_3^1 + 6x_3^{11} &\geq -5 \\
 2x_1 + 3x_2 + 5x_3^1 - 5x_3^{11} &\geq 4 \\
 3x_1 + x_2 + 7x_3^1 - 7x_3^{11} &\geq 3 \\
 -3x_1 - x_2 - 7x_3^1 + 7x_3^{11} &\geq -3
 \end{aligned}$$

$$\text{and } x_1, x_2, x_3^1, x_3^{11} \geq 0$$

$$\begin{aligned}
 & x_1 + 4x_2 + 6x_3 \leq 5 \\
 & 2x_1 + 3x_2 + 5x_3 \geq 4 \\
 & 3x_1 + x_2 + 7x_3 = 3 \\
 & \text{and } x_1, x_2 \geq 0, \\
 & x_3 \text{ is unrestricted.}
 \end{aligned}
 \quad \left| \begin{array}{l}
 -x_1 - 4x_2 - 6(x_3^I - x_3^{II}) \geq -5 \\
 2x_1 + 3x_2 + 5(x_3^I - x_3^{II}) \geq 4 \\
 3x_1 + x_2 + 7(x_3^I - x_3^{II}) \geq 3 \\
 \hline
 3x_1 + x_2 + 7(x_3^I - x_3^{II}) \leq 3 \\
 -3x_1 - x_2 - 7(x_3^I - x_3^{II}) \geq -3
 \end{array} \right. \times -1$$

Standard Primal

$$\text{Min } Z = 3x_1 + 5x_2 + 6x_3^I - 6x_3^{II}$$

s.t.

$$-x_1 - 4x_2 - 6x_3^I + 6x_3^{II} \geq -5$$

$$2x_1 + 3x_2 + 5x_3^I - 5x_3^{II} \geq 4$$

$$3x_1 + x_2 + 7x_3^I - 7x_3^{II} \geq 3$$

$$-3x_1 - x_2 - 7x_3^I + 7x_3^{II} \geq -3$$

$$\text{and } x_1, x_2, x_3^I, x_3^{II} \geq 0$$

$$\text{Min } Z = \underline{\underline{3x_1 + 5x_2 + 6x_3^1 - 6x_3^{11}}}$$

s.t.

$$\begin{aligned} -x_1 - 4x_2 - 6x_3^1 + 6x_3^{11} &\geq -5 \\ 2x_1 + 3x_2 + 5x_3^1 - 5x_3^{11} &\geq 4 \\ 3x_1 + x_2 + 7x_3^1 - 7x_3^{11} &\geq 3 \\ -3x_1 - x_2 - 7x_3^1 + 7x_3^{11} &\geq -3 \end{aligned}$$

and  $x_1, x_2, x_3^1, x_3^{11} \geq 0$

Dual

$$\text{Max } Z = -5w_1 + 4w_2 + 3w_3 - 3w_4$$

s.t.

$$\begin{aligned} -w_1 + 2w_2 + 3w_3 - 3w_4 &\leq 3 \\ -4w_1 + 3w_2 + w_3 - w_4 &\leq 5 \\ -6w_1 + 5w_2 + 7w_3 - 7w_4 &\leq 6 \\ 6w_1 - 5w_2 - 7w_3 + 7w_4 &\leq -6 \end{aligned}$$

and  $w_1, w_2, w_3, w_4 \geq 0$

$$\text{Min } Z = \boxed{3x_1 + 5x_2 + 6x_3^1 - 6x_3^{11}}$$

s.t.

$$-x_1 - 4x_2 - 6x_3^1 + 6x_3^{11} \geq \boxed{-5}$$

$$2x_1 + 3x_2 + 5x_3^1 - 5x_3^{11} \geq \boxed{4}$$

$$3x_1 + x_2 + 7x_3^1 - 7x_3^{11} \geq \boxed{3}$$

$$-3x_1 - x_2 - 7x_3^1 + 7x_3^{11} \geq \boxed{-3}$$

$$\text{and } x_1, x_2, x_3^1, x_3^{11} \geq 0$$

Dual

$$\text{Max } Z = -5w_1 + 4w_2 + 3w_3 - 3w_4$$

s.t.

$$-w_1 + 2w_2 + 3w_3 - 3w_4 \leq \boxed{3}$$

$$-4w_1 + 3w_2 + w_3 - w_4 \leq \boxed{5}$$

$$-6w_1 + 5w_2 + 7w_3 - 7w_4 \leq \boxed{6}$$

$$6w_1 - 5w_2 - 7w_3 + 7w_4 \leq \boxed{-6}$$

$$\text{and } w_1, w_2, w_3, w_4 \geq 0$$

$\tilde{w} =$   
 $\underline{(w_3 - w_4)}$