

# Numerical Methods I

## Exam #1 Tips

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# Exam #1

## 0.1 Absolute Error

$$e_{abs} = |e_{exact} - e_{approx}| \quad (1)$$

## 0.2 Relative Error

$$e_{rel} = \frac{|e_{exact} - e_{approx}|}{|e_{exact}|} \quad (2)$$

## 0.3 Bisection Method

$$m = \frac{b + a}{2} \quad (3)$$

If  $y < 0$ ,  $a = m$ ;  
Else if  $y > 0$ ,  $b = m$ ;  
Else  $y = 0$ .

## 0.4 Regula Falsi

$$x = b - y_b \frac{b - a}{y_b - y_a} \quad (4)$$

If  $y < 0$ ,  $a = x$ ;  
Else if  $y > 0$ ,  $b = x$ ;  
Else  $y = 0$ .

## 0.5 Secant Method

$$x_{n+1} = x_n - \frac{f(x_n) * (x_n - x_{n-1})}{f(x_n) - f(x_{n-1})} \quad (5)$$

## 0.6 Newton's Method

$$x_{n+1} = x_n - \frac{y(x_n)}{y'(x_n)} \quad (6)$$

## 0.7 Fixed-Point

$$f(p) = p \quad (7)$$

$$x = f(x) = 0 \quad (8)$$

Find a convenient  $x = g(x)$ , and use:

$$x_{n+1} = g(x_n) \quad (9)$$

## 0.8 Existence (Intermediate Value Theorem)

If  $y$  changes signal in  $[a, b]$  and is continuous, exists at least one root in  $[a, b]$ .

## 0.9 Uniqueness (Monotone)

If  $y'$  does not change signal in  $[a, b]$  and it is continuous, exists only one root in  $[a, b]$ .

## 0.10 Newton's First Modification Method

$$x_{n+1} = x_n - \frac{my(x_n)}{y'(x_n)} \quad (10)$$

Used when multiplicity of the root is known to be  $m > 1$ . Converges Quadratically.

## 0.11 Newton's Second Modification Method

$$x_{n+1} = x_n - \frac{y(x_n)y'(x_n)}{[y'(x_n)]^2 - y(x_n)y''(x_n)} \quad (11)$$

Used when multiplicity of the root is not known. Converges Quadratically.

## 0.12 Discovering Multiplicity

If  $x = \alpha$  is root of the function. Do continuous derivatives until:

$$f(a) = 0$$

$$f'(a) = 0$$

$$f^{(n-1)}(a) = 0$$

$$f^{(n)}(a) \neq 0 \quad (12)$$

$$(13)$$

The function in root  $x = \alpha$  have multiplicity  $n$ .

## 0.13 Order Convergence

$$|e_{n+1}| \approx \beta |e_n|^R \quad (14)$$

$\beta$  is the asymptotic error constant.  $R$  is the order of convergence.

## 0.14 Fixed-Point Theorem

If  $g(x)$  is continuous on  $[a, b]$  and  $g(x) \in [a, b] \forall x \in [a, b]$ . Then,  $g$  is guaranteed to have at least one fixed-point in  $[a, b]$ .

Moreover, if  $|g'(x)| \leq k < 1 \forall x \in [a, b]$ , then the fixed-point is unique.

Step one: Verify the boundaries. If positive,  $g$  have at least one fixed-point. If negative,  $g$  does not have fixed points. Don't forget to check if  $g$  is monotone.

$$g(a) \in [a, b]? \quad (15)$$

$$g(b) \in [a, b]? \quad (16)$$

$$g'(x) \geq 0 \forall x \in [a, b]?, \text{ or} \quad (17)$$

$$g'(x) \leq 0 \forall x \in [a, b]? \quad (17)$$

Step two: Verify if the modular derivated function have maximum value less than 1. If  $|g'(x)| > 1$ , the sequence will diverge. Otherwise, the fixed point is unique and can be discovered by the iterative scheme.

$$x_{n+1} = g(x_n), \forall x_0 \in [a, b] \quad (18)$$

## 0.15 Bisection Error Bound

$$|E| \leq \frac{b-a}{2^n} \quad (19)$$

To discover the minimum number of iterations needed to ensure a required tolerance ( $T$ ):

$$\lceil n \rceil \geq \frac{\ln \frac{b-a}{T}}{\ln 2} \quad (20)$$

## 0.16 Fixed-Point Error Bound

$$|E| \leq \frac{k^n}{1-k} |x_0 - x_1| \quad (21)$$

To discover the minimum number of iterations needed to ensure a required tolerance ( $T$ ):

$$\lceil n \rceil \geq \frac{\ln \frac{T(1-k)}{|x_0-x_1|}}{\ln k} \quad (22)$$

To discover the asymptotic error constant for an  $x = \alpha$  as fixed point of  $g$ .

$$\lim_{n \rightarrow \infty} \frac{|e_{n+1}|}{|e_n|} = \frac{g^{(R)}(\alpha)}{R!} \quad (23)$$