# Problem Set V

ECO 7427 - Econometric Theory II

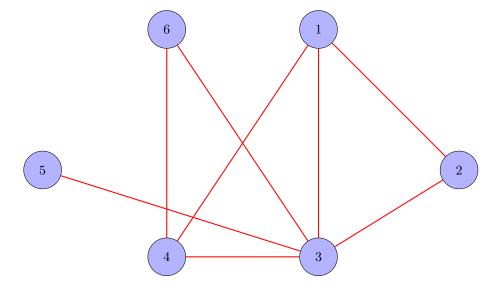
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# Social Networks

Let  $N = \Big\{1, 2, 3, 4, 5, 6\Big\}$  and consider the following network (N, g).

Figure 1: Undirected and Unweighted Network



#### Construct the adjacency matrix:

$$g = \begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 \end{bmatrix}$$

To compute the degree of each vertex we sum across each row

$$d_i(g) = \sum_j g_{ij} \tag{1}$$

## The neighborhood of each vertex:

Table 1: Neighborhood of each vertex

| Vertex | $N_i(g)$            |
|--------|---------------------|
| N = 1  | $\{2, 3, 4\}$       |
| N = 2  | $\{1,3\}$           |
| N = 3  | $\{1, 2, 4, 5, 6\}$ |
| N = 4  | $\{1, 3, 6\}$       |
| N = 5  | $\{3\}$             |
| N = 6  | $\{3,4\}$           |

#### The degree of each node:

Table 2: Degree of each vertex in the network

| Vertex | Degree |
|--------|--------|
| N = 1  | 3      |
| N = 2  | 2      |
| N = 3  | 5      |
| N = 4  | 3      |
| N = 5  | 1      |
| N = 6  | 2      |

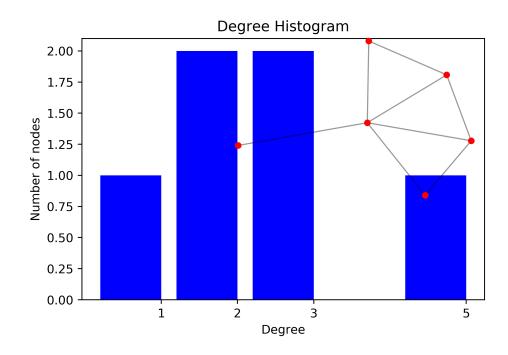
Average degree of the network  $g \in G(N)$ :

$$\frac{1}{N} \sum_{i \in N} d_i(g) = \frac{1}{6} \left[ 3 + 2 + 5 + 3 + 1 + 2 \right] = \frac{8}{3}$$
 (2)

The density of the network  $g \in G(N)$ :

$$\frac{\sum_{i \in N} d_i(g)}{n(n-1)} = \frac{\frac{1}{n} \sum_{i \in N} d_i(g)}{n-1} = \frac{8/3}{5} = \frac{8}{15}$$
 (3)

Figure 2: Degree Histogram



The individual clustering of the six vertices:

$$CI_{i}(g) = \frac{\left\{ jk \in g | k \neq j, j \in N_{i}(g), k \in N_{i}(g) \right\}}{d_{i}(g) [d_{i}(g) - 1]}$$
(4)

Table 3: Individual clustering of each vertex

| Vertex | $CI_i(g)$ |
|--------|-----------|
| N = 1  | 2/3       |
| N = 2  | 1         |
| N = 3  | 3/10      |
| N = 4  | 2/3       |
| N = 5  | 0         |
| N = 6  | 1         |
| 11 0   | -         |

The average clustering coefficient:

$$CI^{\text{Avg}}(g) = \frac{1}{n} \sum_{i \in N} CI_i(g) = \frac{1}{6} \left[ \frac{2}{3} + 1 + \frac{3}{10} + \frac{2}{3} + 0 + 1 \right] = \frac{109}{180}$$
 (5)

#### Degree centrality for each vertex:

$$CI_i(g) = \frac{\left\{ jk \in g | k \neq j, j \in N_i(g), k \in N_i(g) \right\}}{d_i(g) \left[ d_i(g) - 1 \right]} \tag{6}$$

Table 4: Degree centrality for each vertex

| Vertex | $CI_i(g)$ |
|--------|-----------|
| N = 1  | 3/5       |
| N = 2  | 2/5       |
| N = 3  | 1         |
| N = 4  | 3/5       |
| N = 5  | 1/5       |
| N = 6  | 2/5       |

Hence from Table 4 we can see that vertex 3 is well connected in terms of direct connections as it has the highest degree of centrality. However using this measure alone we cannot tell if vertex 3 is well located in the network (N, g).

#### Closeness centrality fo each vertex:

$$C_i(g) = \frac{1}{\frac{1}{n-1} \sum_{j \neq i} l(i,j)}$$
 (7)

where l(i, j) denotes the number of links in the **shortest path** between i and j.

Table 5: Closeness centrality fo each vertex

| Vertex | $C_i(g)$ |
|--------|----------|
| N = 1  | 5/7      |
| N = 2  | 5/8      |
| N = 3  | 1        |
| N = 4  | 5/7      |
| N = 5  | 5/9      |
| N = 6  | 5/8      |

#### Decay centrality for each vertex:

$$D_i^{\delta}(g) = \sum_{j \neq i} \delta^{l(i,j)} \tag{8}$$

where l(i, j) denotes the number of links in the **shortest path** between i and j.

Table 6: Decay centrality fo each vertex

| Vertex | $D_i^{\delta}(g)$   | $\delta = 1/2$ |
|--------|---------------------|----------------|
| N = 1  | $\delta(3+2\delta)$ | 2              |
| N = 2  | $\delta(2+3\delta)$ | 7/4            |
| N = 3  | $5\delta$           | 5/2            |
| N = 4  | $\delta(3+2\delta)$ | 2              |
| N = 5  | $\delta(1+4\delta)$ | 3/2            |
| N = 6  | $\delta(1+4\delta)$ | 3/2            |

#### Betweenness centrality for each vertex:

$$Ce_i^B(g) = \sum_{k \neq j: i \notin \{k, j\}} \frac{P_i(kj)/P(kj)}{\frac{1}{2}(n-1)(n-2)}$$
(9)

Table 7: Betweenness centrality of a vertex

| Vertex | $P_i(k,j)$                                   | $Ce_i^B(g)$ |
|--------|--|-------------|
| N = 1  | $\{23, 24, 25, 26, 34, 35, 36, 45, 46, 56\}$ | 1/20        |
| N = 2  | $\{13, 14, 15, 16, 34, 35, 36, 45, 46, 56\}$ | 0           |
| N = 3  | $\{12, 14, 15, 16, 24, 25, 26, 45, 46, 56\}$ | 3/5         |
| N = 4  | $\{12, 13, 15, 16, 23, 25, 26, 35, 36, 56\}$ | 1/20        |
| N = 5  | $\{12, 13, 14, 16, 23, 24, 26, 34, 36, 46\}$ | 0           |
| N = 6  | $\{12, 13, 14, 15, 23, 24, 25, 34, 35, 45\}$ | 0           |

Notice that for the first vertex (N=1) the link 24 shown in red in the set  $P_i(k,j)$  can either go through vertex 1 or through vertex 3 since both are geodesic.

#### Eigenvector centrality for each node:

We solve the following problem numerically to compute the eigencentrality of each vertex in the graph.

$$\lambda C^e(g) = gC^e(g) \tag{10}$$

where  $C^e(g)$  is left-hand eigenvector of the adjacency matrix g and  $\lambda$  is the corresponding eigenvalue.

where  $|\lambda_i| \leq \lambda_1 = 3.04 \ (\forall i \neq 1)$  is the **Perron-Frobenius eigenvalue** of the adjacency matrix g also known as the **spectral radius**<sup>1</sup>  $\rho(g)$ . Eigencentrality is a measure of the influence of a node in a network.

Theorem: Let  $g \in C^{n \times n}$  with spectral radius  $\rho(g)$ , then  $\rho_A < 1$  if and only if  $\lim_{g \to \infty} g^k = 0$  and if  $\rho(g) > 1$  then  $\lim_{k \to \infty} \|g^k\| = \infty$ 

Table 8: Largest eigenvalue and its eigenvector

| Vertices | $\lambda_{max} = \rho(g) = 3.04$ |
|----------|----------------------------------|
| N = 1    | 0.45                             |
| N = 2    | 0.34                             |
| N = 3    | 0.58                             |
| N = 4    | 0.45                             |
| N = 5    | 0.19                             |
| N = 6    | 0.34                             |

### Katz-Bonacich Centrality:

$$C^{B}(g, a, b) = (I_{n} - bg)^{-1} ag \mathbf{1}_{n}$$
(11)

Table 9: Number of walks emanating from each vertex

| Vertices | k = 1 | <br>$k = \infty$ |
|----------|-------|------------------|
| N = 1    | 3     | <br>$\infty$     |
| N = 2    | 2     | <br>$\infty$     |
| N = 3    | 5     | <br>$\infty$     |
| N = 4    | 3     | <br>$\infty$     |
| N = 5    | 1     | <br>$\infty$     |
| N = 6    | 2     | <br>$\infty$     |

where we would expect that the number of walks goes to infinity as  $\lim_{k\to\infty}\left\|A^k\right\|=\infty$ 

We compute Katz-Bonacich centrality using an algorithmic approach in Python

$$x_i = \beta \sum_j A_{ij} x_j + \alpha \tag{12}$$

where the parameter  $\alpha$  controls the initial centrality and  $\beta < \frac{1}{\lambda_{\max}}.$ 

Table 10: Katz-Bonacich centrality

| Vertices | $\alpha = 1,  \beta = 0.5$ | $\alpha = 1,  \beta = 1/3$ |
|----------|----------------------------|----------------------------|
| N = 1    | 0.47                       | 0.45                       |
| N = 2    | 0.31                       | 0.34                       |
| N = 3    | 0.60                       | <b>0.58</b>                |
| N = 4    | 0.47                       | 0.45                       |
| N = 5    | 0.08                       | 0.19                       |
| N = 6    | 0.31                       | 0.34                       |

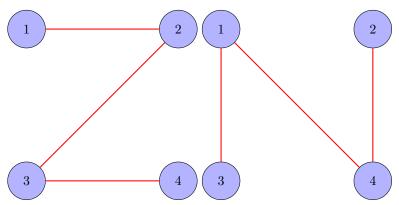
# **Graph Theory**

Show that if a network is not connected, then its complement is.

Proof. Let G = (V, E) be a graph comprising of a set of vertices V and a set of edges E. Let  $\langle v_i, v_j \rangle$  denote an edge between two arbitrary vertices  $v_i$  and  $v_j$  in the graph G where  $\langle v_i, v_j \rangle \in E$  and let  $\tilde{G} = (V, \tilde{E})$  be the complement of the graph G. Since G is not a connected graph then we can partition the graph G into two disjoint sets of vertices  $V_1$  and  $V_2$  where  $\forall v_1 \in V_1$  and  $\forall v_2 \in V_2$  we have that  $\langle v_1, v_2 \rangle \notin E$ . Hence for all  $v_1 \in V_1$  and  $v_2 \in V_2$  we have that  $\langle v_1, v_2 \rangle \in \tilde{E}$  which implies that  $\tilde{G}$  is a connected graph.  $\square$ 

Provide an example of a four-node network that is connected and its complement is also connected:

Figure 3: Left: ({1,2,3,4},g), right: ({1,2,3,4},g')



# **Correlation Scatter Plots**

Figure 4: Scatter plot

## Simple Scatterplot Matrix

