

## 7. Higher Order Operations on Lists

[ Reade §3.4 ]

### Examples.

(1) To *double* the values of a list of integers:

e.g. `double_list [3, 1, 4] = [6, 2, 8]`.

Def. by *list recursion*:

```
fun double_list (a::x) = (2 * a)::(double_list x)
  | double_list [ ] = [ ];
```

Type?

- We want a *more general* (higher order) operation on lists, from which we can derive `double_list` as a special case!

(2)  $\text{map } f [a_1, \dots, a_n] = [f a_1, \dots, f a_n]$  (“Map  $f$  over list”)

```
fun map f [ ] = [ ]
  | map f (a::x) = (f a)::(map f x);
```

Type?

Then we can define

```
fun double x = 2 * x;
val double_list =
```

(3) To **filter** out of a list  $\ell$  those items not satisfying a predicate  $p$ :

i.e.,  $\text{filter } p \ell = \text{sublist of } \ell \text{ containing those items which satisfy } p$ .

```
fun filter p [ ] = [ ]
  | filter p (a::x) = if p a
                      then a::(filter p x)
                      else filter p x;
```

Definition by *list recursion* on *2nd argument*.

Type?

(4) **Quantifiers** (*Transform predicates on  $T$  to predicates on  $T$  list*):

(a)  $\text{all } p \ell = \begin{cases} \text{true} & \text{if } \forall a \text{ in } \ell : p a \\ \text{false} & \text{otherwise} \end{cases} \quad (p \text{ is a predicate})$

```
fun all p (a::x) = (p a) andalso (all p x)
  | all p [ ]     = ...
```

Def. is by ... ?

Type?

(b)  $\text{exists } p \ell :$  **Ex**

(5)  $\text{addlists}([a_1, \dots, a_n], [b_1, \dots, b_n]) = [a_1 + b_1, \dots, a_n + b_n]$

```
fun addlists ([ ], [ ]) = [ ]
  | addlists (a::x, b::y) = (a + b)::(addlists (x, y))
```

*! Warning: pattern matching is not exhaustive*

Type?

### Notes.

- (1) Definition by *list recursion* on *both arguments*.
- (2) If arguments have different lengths: **error!** (No pattern matching.)

- Again, we want a *more general* (higher order) operator:

(6)  $\text{zip } f ([a_1, \dots, a_n], [b_1, \dots, b_n]) = [f(a_1, b_1), \dots, f(a_n, b_n)]$

```
- fun zip f (a::x, b::y) = (f(a,b))::(zip f (x, y))
  | zip f ([ ], [ ]) = [ ];
```

*! Warning: pattern matching is not exhaustive*

```
- zip plus ([2], [3,5]);
```

*! Uncaught exception*

*! Match*

Type?

Then we can define

```
val addlists = zip op+;
val multlists = zip op*;
```

etc.

**Note:** Again, get *error* with lists of different lengths.

**Ex** Define variants of `zip` which work with lists of different lengths:

- (a) Append tail of longer list,
- (b) Chop off tail of longer list.

(7) `quicksort : int list → int list` [See Reade p. 108]

Very nice! — it uses higher order operators (including `filter`).

(8) To *sum* a list of integers:

E.g. `sumlist [2, 6, 3] = 11`.

Two algorithms:

(i) *Primitive (list) recursion*:

```
fun sumlist [ ] = 0
  | sumlist (b::x) = b + (sumlist x);
```

Type?

(ii) *Tail (list) recursion*:

```
local fun sumlist_iter (a, [ ]) = a
      | sumlist_iter (a, b::x) = sumlist_iter (a + b, x)
in fun  sumlist x = sumlist_iter (0, x)
end;
```

### Notes.

(1) Def. is by *list recursion* on *2nd argt.*

Note the *extra argument* `a` in `sumlist_iter`.

(Think of `a` as the “partial sum”, which *grows* from 0 to `sumlist`, and `x` as the “partial list”, which *shrinks* from the input list to `[ ]`.)

(2) Iteration invariant = `sumlist_iter (a, x) = ... ?`

Bound value = ... ? **Ex**

**Ex** Define `prodlist`  $\ell$  = product of list  $\ell$  of integers, analogous to (i) and (ii) above.

Again, we want a *more general (higher order)* operator.

We will generalise each of the above two definitions of `sumlist`.

(9)(i) Define `reduce f a  $\ell$`  where

$f$  is a binary function

$a$  is *starting value* (for empty list)

$\ell$  is a list.

```
fun reduce f a [ ] = a
  | reduce f a (b::x) = f (b, (reduce f a x));
```

**Note:** Definition is by (*prim.*) *list recursion* on *3rd argt.*

Type?

Hence:

- (a) `val sumlist = reduce ...`
- (b) `val prodlist = reduce ...`
- (c) Let `flatten` be a function such that (e.g.)

`flatten [[1, 2], [3], [4, 5, 6]] = [1, 2, 3, 4, 5, 6]`

**Ex** `val flatten =`  
Type?

**Ex** Define `minlist  $\ell$  = minimum of a list  $\ell$  of integers.`  
( If  $\ell$  is empty, return an error message.)

- (9)(ii) **Accumulate:** `accum f a  $\ell$`  is a *tail recursive* version of `reduce`.  
(Again, think of  $a$  as a *partial result* which grows, and  $\ell$  as a *partial list* which shrinks.)

`fun accum f a [ ] = a`  
| `accum f a (b::x) = accum f (f (a, b)) x;`

Type?

**Note:** We can define

`sumlist,`  
`prodlist,`  
`flatten,`  
`minlist, ...`

from `accum` exactly as from `reduce` in (i), e.g.:

`val sumlist = accum (op +) 0;`

Informally *compare* these two functionals (‘ $\otimes$ ’ is any binary operator in infix):

$$(i) \text{ reduce } (\text{op } \otimes) a [b_1, b_2, b_3] = b_1 \otimes (b_2 \otimes (b_3 \otimes a))$$

$$(ii) \text{ accum } (\text{op } \otimes) a [b_1, b_2, b_3] = ((a \otimes b_1) \otimes b_2) \otimes b_3$$

‘reduce’ is also called ‘foldright’,

‘accum’ is also called ‘foldleft’.

**Note.** ‘op  $\otimes$ ’ is the *prefix* (uncurried) form of ‘ $\otimes$ ’.

**Ex** Give reasonable *sufficient conditions* on  $f$  and  $a$  such that for any list  $\ell$

$$\text{reduce } f a \ell = \text{accum } f a \ell.$$