

Kae-Yang Hsieh(001092745)

Program Structures & Algorithms

Fall 2021 Assignment No. 2

Tasks

*** Observations of the insertion sort regarding the order of growth**

Conclusion:

There are four different initial array ordering type: random, ordered, partially ordered, and reverse ordered and each array were be benchmarked with 10 different sized elements. With the removal of the first 3 smaller elements, the four benchmark experiments matched the expectation with the time complexity $O(n)$ of ordered array and $O(n^2)$ of the remaining types. Also, with the same elements, the order of time consumed by the four type were: ordered < partially ordered < random < reverse ordered, which is also matched the expectation.

Preparations

The four different initial arrays were created by the following methods:

- * random: randomly shuffled an ordered array
- * ordered: an array start from 1 to n
- * partially ordered: random shuffled half of an ordered array, start from index 1 to n/2
- * reverse ordered: an array start from n to 1

Each type of the array was initialized with 11 different lengths, followed the doubling method, where n starts from 50 to 51200

The ordered array where n starts from 500 to 512000

Also, the first 50 element is for warmup and will not be accounted into the experiments.

Experiments

The experiments were implemented with a class called InsertionSortBenchmarks, located at "src/test/java/edu/neu/coe/info6205/util/InsertionSortBenchmarks.java"

All the arrays were set up using @before methods and timer.repeat to calculate the mean time of the sort.

Benchmark Results and Analysis

Random array, partially ordered array, and reverse ordered array:

The average case (random and partially ordered) and the worst case (reverse ordered) of the insertion sort will be $O(n^2)$, which was expected to fit with the statistics. However, the equation showed a large difference with the expectation (Figure 1-1, 1-2, 1-3). With the removal of the first four abnormal statistics, the graph (Figure 2-1, 2-2, 2-3) is more close to the expectation: $y = n^2$

Through this experiments, the elements under 400 seems like has more deviation with the expectation. Here is a theory that may cause this situation: the cpu time contained the cache reading time and processing time, the cache reading will be almost the same while the data are smaller than the bandwidth. Also, the cache reading time may have a large proportion when measuring the sorting of little elements. This may affect a lot while we want to measure the processing time but not the cache reading time.

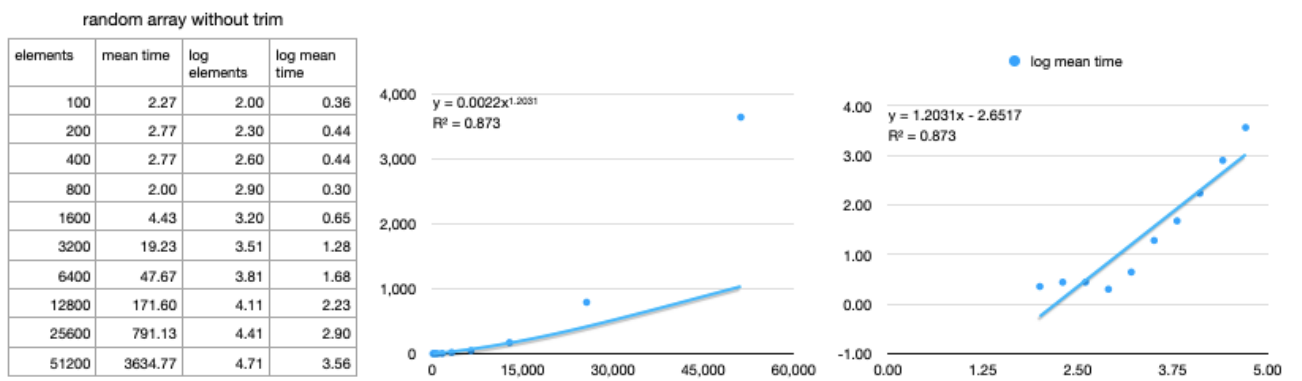


Figure 1-1 Random array without trimming

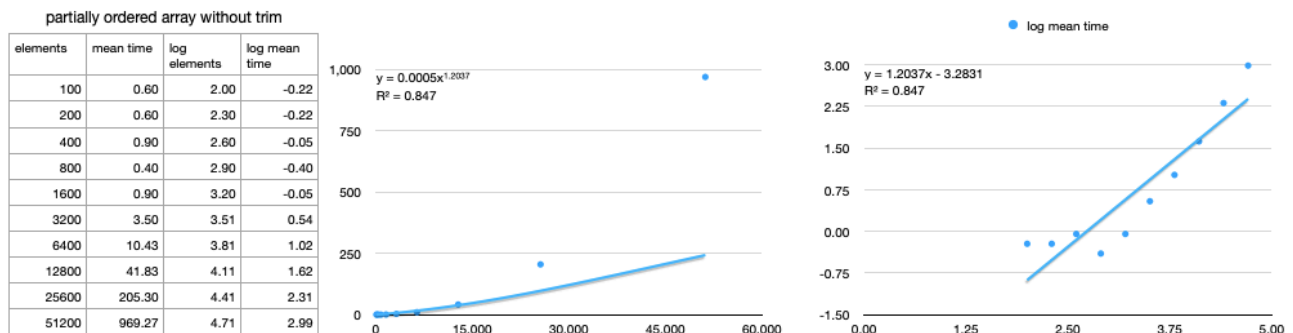


Figure 1-2 Partially ordered array without trimming

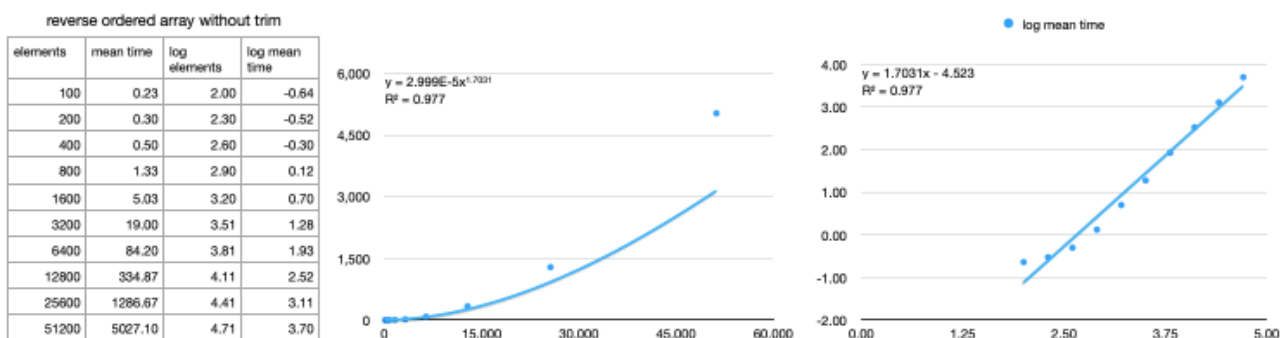


Figure 1-3 Reverse ordered array without trimming

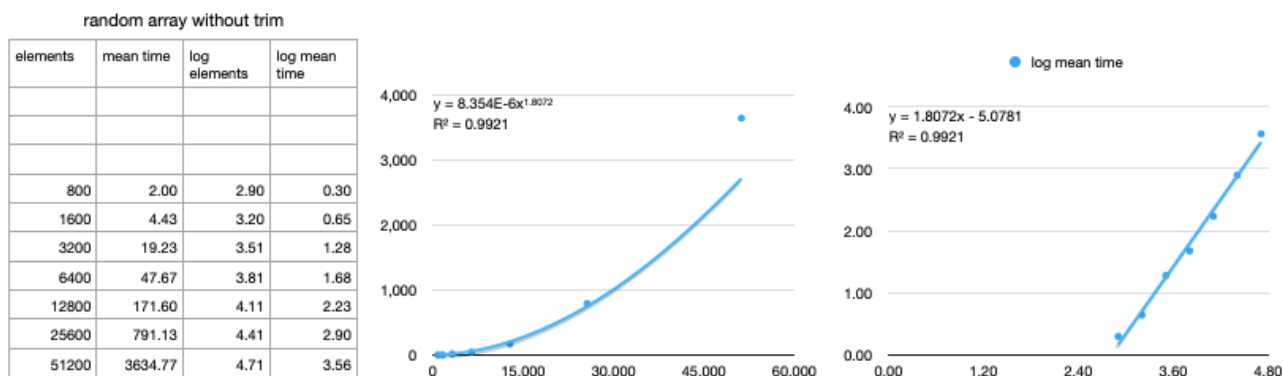


Figure 2-1 Random array with trimming

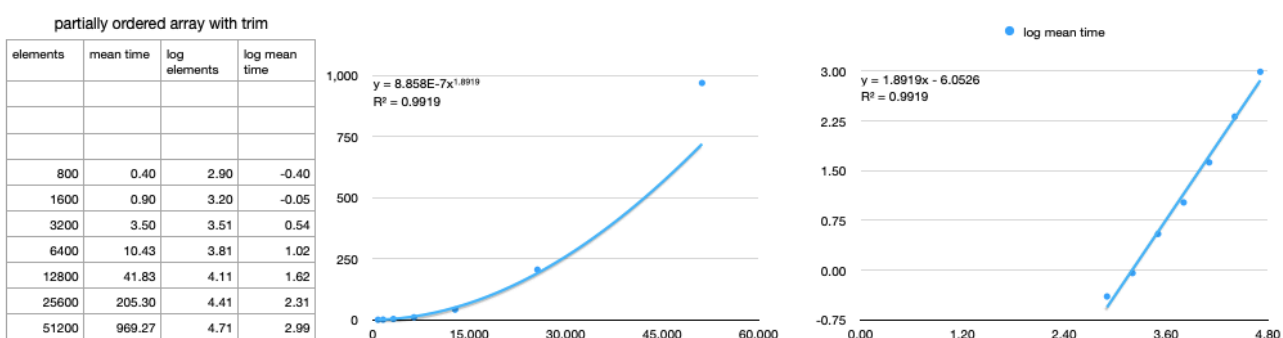


Figure 2-2 Partially ordered array with trimming

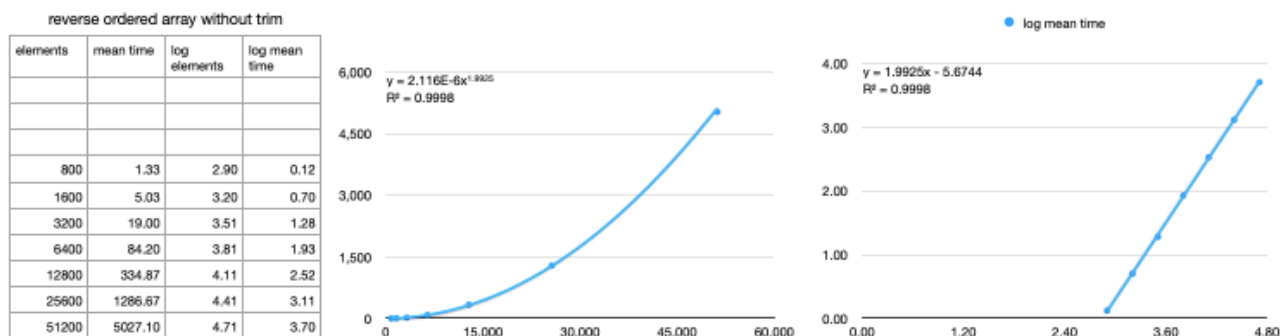


Figure 2-3 Reverse ordered array with trimming

Ordered array:

The ordered array matched the best case of the insertion sort, which has $O(n)$ time complexity. In the ordered array experiments, the original element number was too small to exclude the deviations. The statistic showed the original numbers has no linearity(Figure 3-1). With the ten times of the elements(Figure 3-2), the R^2 is near 1, which is expected to show a $O(n)$ time complexity.

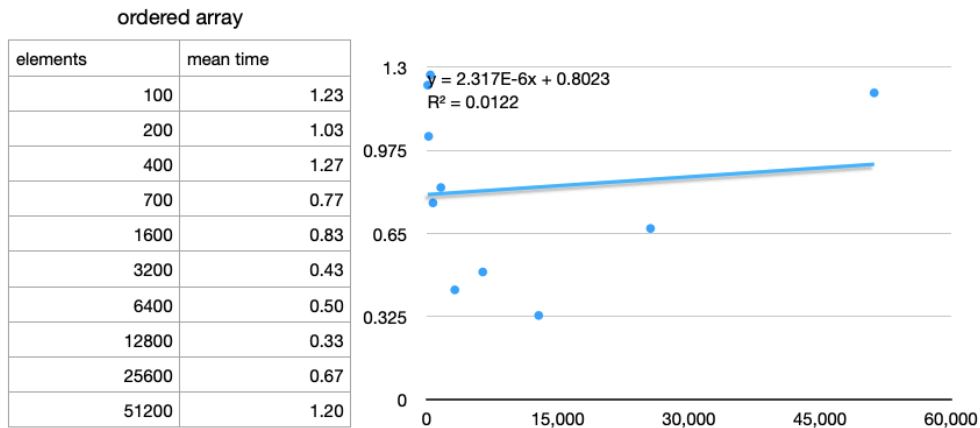


Figure 3-1 Ordered array with original numbers

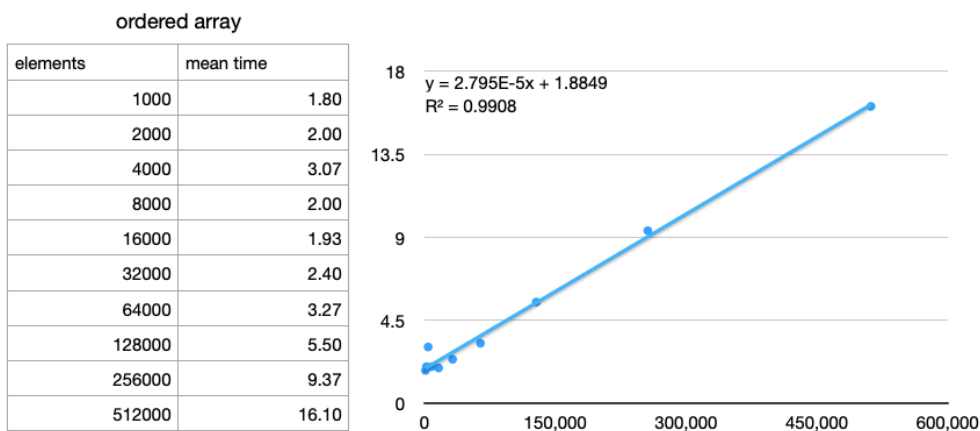
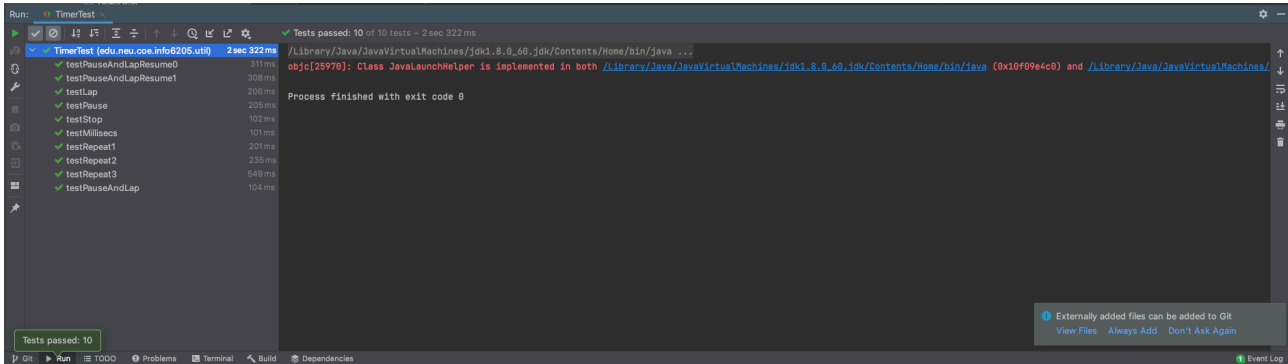


Figure 3-2 Ordered array with ten times numbers

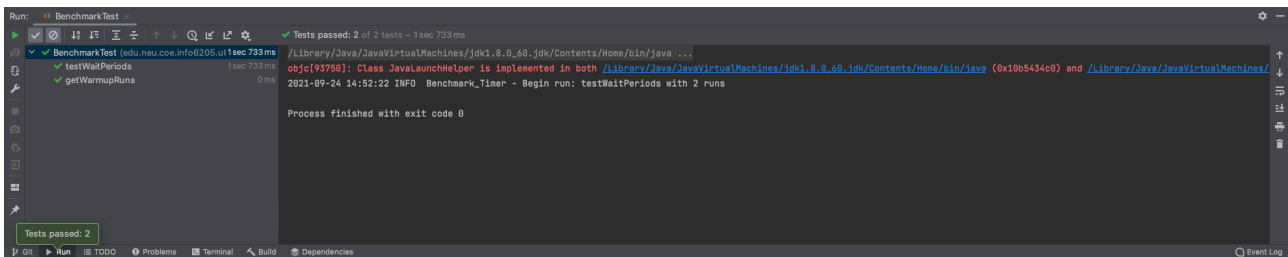
* Unit test results

TimerTest

Note: The delta of the testRepeat1 has been adjusted from 6 to 13 to fit my environment.



BenchmarkTest



InsertionSortTest

