

$$\langle x, x \rangle \geq 0$$

$$\langle x, x \rangle = 0 \Leftrightarrow P(x=0) = 1$$

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• V non-null vector space
with a norm $\|\cdot\|$

Define $d: V \times V \rightarrow [0, \infty)$

$$\text{by } d(x, y) = \begin{cases} \|x\| + \|y\| & \text{if } x \neq y \\ 0 & \text{if } x = y \end{cases}$$

Ex. show d is a metric on V . s.t. \nexists a norm on V that induces d .

• Recall ~~x, y~~ X, Y are independent if

$$P(X=x, Y=y) = P(X=x) P(Y=y)$$

for all ~~$x \in X$ & $y \in Y$~~

$x \in D_X$ & $y \in D_Y$.

(Equivalently,

$$P(X \in B_1, Y \in B_2) = P(X \in B_1) P(Y \in B_2)$$

for any $B_1, B_2 \in \mathcal{R}$)

• Defn. X_1, \dots, X_n are mutually independent

$$\text{if } P(X_1 = x_1, \dots, X_n = x_n)$$

$$= P(X_1 = x_1) \dots P(X_n = x_n)$$

$$\forall x_1, x_2, \dots, x_n.$$

this is equiv. to $P(X_1 \in B_1, \dots, X_n \in B_n)$

$$= P(X_1 \in B_1) \dots P(X_n \in B_n)$$

for any choice of y

$$1 \leq t_1 < \dots < t_K \leq n.$$

$$P(X_{t_1} = y_1, \dots, X_{t_K} = y_K)$$

$$= P(X_{t_1} = y_1) \dots P(X_{t_K} = y_K)$$

[Note: If we take $B_1 = \{x_1\}$, $B_2 = \{x_2\}$ & all other $B_i = \mathbb{R}$ then we are done!]

⑩ Defn. X_1, \dots, X_n are pairwise indep. if X_i, X_j are indep. for all $i \neq j$

[This is a much weaker condn.

Ex. show 3 r.v.s which are pairwise indep. but not indep.]

⑪ Defn. A family $\{X_\alpha : \alpha \in \Lambda\}$ of discrete r.v.s are said to be mut independent if every finite subcollection is indep.

Examples X_1, X_2, X_3 indep.

claim: For any $B_0 \in \mathbb{R}^2$, $C \in \mathbb{R}$

$$P((X_1, X_2) \in B, X_3 \in C)$$

$$= P((X_1, X_2) \in B) P(X_3 \in C)$$

$$\text{LHS} = \sum_{\substack{x_1, x_2, x_3 \\ (x_1, x_2) \in B \\ x_3 \in C}} P(X_1 = x_1, X_2 = x_2, X_3 = x_3)$$

$$\begin{aligned}
&= \sum P(X_1=x_1) P(X_2=x_2) P(X_3=x_3) \\
&= \sum P(X_1=x_1, X_2=x_2) P(X_3=x_3) \\
&= \sum P(X_1=x_1, X_2=x_2) \sum P(X_3=x_3) \\
&= P((X_1, X_2) \in B) P(X_3 \in C)
\end{aligned}$$

$$\begin{aligned}
&V(X_1 + \dots + X_n) \\
&= \sum_i V(X_i) + 2 \sum_{i \neq j} \text{cov}(X_i, X_j)
\end{aligned}$$

$$\begin{aligned}
\text{If } LHS &= \text{cov}(\sum X_i, \sum X_i) \\
&= \dots
\end{aligned}$$

$$\begin{aligned}
&V(X_1 + \dots + X_n) = V(X_1) + \dots + V(X_n) \\
&\text{if } X_1, \dots, X_n \text{ are pairwise independent.}
\end{aligned}$$

Ex. Each purchase of a product gives you one of r different toys

X = # diff. types of toys after n choices

Y = # purchases needed to get all r types.

$$\textcircled{1} X = I_1 + \dots + I_r$$

where $I_j = \begin{cases} 1 & \text{if } j\text{-th type is obtained} \\ 0 & \text{if not.} \end{cases}$

$$E(I_j) = P(j\text{-th type is there})$$

$$= 1 - \left(\frac{r-1}{r}\right)^n$$

$$\therefore E(X) = r \left(1 - \left(\frac{r-1}{r}\right)^n\right)$$

$$V(X) = \sum_{j=1}^r V(I_j) + \sum_{j < k} \text{cov}(I_j, I_k)$$

$$\begin{aligned} V(I_j) &= E(I_j) - [E(I_j)]^2 \\ &= E(I_j) (1 - E(I_j)) \\ &= \left(\frac{r-1}{r}\right)^n \left[1 - \left(\frac{r-1}{r}\right)^n\right] \end{aligned}$$

$$\begin{aligned} \text{cov}(I_j, I_k) &= E(I_j I_k) - E(I_j) E(I_k) \\ &= 1 - 2 \left(\frac{r-1}{r}\right)^n + \left(\frac{r-2}{r}\right)^n - \left(1 - \left(\frac{r-1}{r}\right)^n\right)^2 \\ &= \underline{\underline{\left(\frac{r-2}{r}\right)^n - \left(\frac{r-1}{r}\right)^n}} \end{aligned}$$

from inclusion-exclusion

so, finally $V(X) = \dots$

$$Y = Y_1 + Y_2 + Y_3 + \dots + Y_r.$$

clearly

$$Y_1 = 1$$

$$Y_2 \sim \text{Geo}\left(\frac{r-1}{r}\right)$$

$$Y_3 \sim \text{Geo}\left(\frac{r-2}{r}\right)$$

$$\vdots$$

$$Y_j \sim \text{Geo}\left(\frac{r-j+1}{r}\right)$$

$$\vdots$$

$$Y_r \sim \text{Geo}\left(\frac{1}{r}\right)$$

$$E(Y) = \sum E(Y_i)$$

$$= 1 + \frac{r}{r-1} + \frac{r}{r-2} + \dots + \frac{r}{1}$$

$$\left[\dots E(Y_j) = \frac{r}{r-j+1} \right]$$

$$= r \left(\frac{1}{r} + \frac{1}{r-1} + \dots + 1 \right)$$

$$V(Y) = \sum_1^r V(Y_i) + 2 \sum_{j < k} \text{cov}(Y_j, Y_k)$$

$$\textcircled{X} X \sim \text{Geo}(p)$$

$$E(X(X-1)) = \sum_{k=1}^{\infty} k(k-1) p q^{k-1}$$

$$= p q \sum_{k=2}^{\infty} \underbrace{k(k-1) q^{k-2}}_{\frac{d^2}{dq^2}(q^k)}$$

$$= p q \sum \frac{d^2}{dq^2}(q^k)$$

$$= p q \frac{d^2}{dq^2} \sum q^k$$

Easy calculation!

$$\left\{ \begin{aligned} &= p q \left(\frac{d^2}{dq^2} \left(\frac{1}{1-q} \right) \right) = \frac{2q}{p^2} \end{aligned} \right.$$

$$\text{Var}(X) = \frac{2q}{p^2} + \frac{1}{p} - \frac{1}{p^2}$$

$$= \dots$$

$$= \frac{q}{p^2}$$

$$\text{Now, } q = \frac{j-1}{r}$$

$$p = \frac{r-j+1}{r}$$

$$\therefore \sum V(Y_j) = \sum_{j=2}^r \frac{j-1}{r} \frac{r^2}{(r-j+1)^2}$$

Ex. Now, show that the r.v.'s Y_j 's are pairwise indep. so, $\text{cov}(Y_j, Y_k) = 0$.

$$\text{so, } V(Y) = \sum_{j=2}^r \frac{j-1}{r} \frac{r^2}{(r-j+1)^2}$$

X, Y discrete r.v.'s.

$$p(x, y) = P(X=x, Y=y)$$

given $Y=y$,

the conditional dist. of X .

$$P(X|Y) = P(X=x | Y=y)$$

$$= \frac{P(x, y)}{\sum_x P(x, y)}$$

Ex. $X = \#$ signal transmitted from a transmission centre.

$$X \sim \text{Poi}(\lambda)$$

The signal transmitted is received by the reception centre with prob. p & not received by prob $q = 1 - p$; independently of other signals.

$Y = \#$ signals received.

$$D = \{ (i, j) \mid 0 \leq j \leq i \}$$

i, j integers.

$$\begin{aligned} p(i, j) &= P(X=i, Y=j) \\ &= P(Y=j \mid X=i) \underbrace{P(X=i)}_{e^{-\lambda} \frac{\lambda^i}{i!}} \\ &= \binom{i}{j} p^j q^{i-j} \left(e^{-\lambda} \frac{\lambda^i}{i!} \right) \end{aligned}$$

$$\begin{aligned} P(X=i \mid Y=j) &= \frac{e^{-\lambda} \frac{\lambda^i}{i!} \binom{i}{j} p^j q^{i-j}}{\sum_{k=1}^{\infty} e^{-\lambda} \frac{\lambda^k}{k!} \binom{k}{j} p^j q^{k-j}} \end{aligned}$$

Let, $Z = X - Y = \#$ missed signals.

$$\begin{aligned} P(Z=K \mid Y=j) &= \frac{P(Z=K, Y=j)}{P(Y=j)} \\ &= P(X=K+j, Y=j) \\ &= e^{-\lambda} \frac{\lambda^{K+j}}{(K+j)!} \binom{K+j}{j} p^j q^K \end{aligned}$$

$$= e^{-\lambda} \frac{\lambda^{k+j}}{k!j!} p^j q^k$$

~~$$= e^{-\lambda p} \frac{(\lambda p)^j}{j!} \left[e^{-\lambda q} \frac{(\lambda q)^k}{k!} \right]$$~~

$$= \left[e^{-\lambda p} \frac{(\lambda p)^j}{j!} \right] \left[e^{-\lambda q} \frac{(\lambda q)^k}{k!} \right]$$

so, they are independent.

(*) Poisson & Normal dists. can be written as some K numbers of same distribution; i.e. ~~K numbers~~ we can divide it in K poisson or K normal distribution.

Ex. Coin with $p(H) = p$.
 $X = \#$ tossed ^{position} with 7th H
 $Y = \#$ " " 10th H.

$$D = \{ (i, j) \mid 7 \leq i < i+3 \leq j \}$$

$$P(X=i, Y=j) = \binom{i-1}{6} \binom{j-i-1}{2} p^{10} q^{j-10}$$

$$P(X=i \mid Y=j) = \frac{\binom{i-1}{6} \binom{j-i-1}{2} p^{10} q^{j-10}}{\binom{j-1}{9} p^{10} q^{j-10}}$$

(semi) hyper. geo. distribution.

~~Task~~ Toss a coin

$$P(H) = p$$

Times of 1st H, 2nd H & 3rd H.
as a vector in \mathbb{R}^3 .

Let $X_1 \rightarrow$ time till 1st H

$X_2 \rightarrow$ " " 2nd H

$X_3 \rightarrow$ " " 3rd H.

Find cond. prob. of (X_1, X_2) given X_3 .

More generally (X_1, X_2, \dots, X_n) give X_{n+1} .

Analysis

Weierstrass Approx. Thm.

Let, $f: [0,1] \rightarrow \mathbb{R}$ be a
cont. fn.

Fact 1: $\exists M > 0$ s.t.

$$|f(t)| \leq M \quad \forall t \in [0,1]$$

Fact 2, given any $\varepsilon > 0$, $\exists \delta > 0$.

Uniforms cond.

Fix an $\varepsilon > 0$ for any $n \geq 1$

$$P_n(t) = \sum_{k=0}^n \binom{n}{k} f\left(\frac{k}{n}\right) t^k (1-t)^{n-k}$$

$t \in [0,1]$

P_n is a polynomial of degree n .

$$|f(t) - P_n(t)|$$

$$= \left| f(t) - \sum_{k=0}^n \binom{n}{k} f\left(\frac{k}{n}\right) t^k (1-t)^{n-k} \right|$$

$$= \left| \sum_{k=0}^n \left(f(t) - f\left(\frac{k}{n}\right) \right) \binom{n}{k} t^k (1-t)^{n-k} \right|$$

$$\leq \sum_{k=0}^n \left| f(t) - f\left(\frac{k}{n}\right) \right| \binom{n}{k} t^k (1-t)^{n-k}$$

$$= \sum_{\substack{0 \leq k \leq n \\ \left| \frac{k}{n} - t \right| < \delta}} \dots$$

$$\left| \frac{k}{n} - t \right| < \delta$$

$$+ \sum_{\substack{0 \leq k \leq n \\ \left| \frac{k}{n} - t \right| \geq \delta}} \dots$$

$$\left| \frac{k}{n} - t \right| \geq \delta$$

Now the first sum

$$< \sum_{0 \leq k \leq n} \binom{n}{k} t^k (1-t)^{n-k}$$

this term is
bdd. by 1.

$$\text{2nd sum} \leq 2M \sum_{k: \left| \frac{k}{n} - t \right| \geq \delta} \binom{n}{k} t^k (1-t)^{n-k}$$

If $X \sim \text{Bin}(n, t)$ then,

$$\text{2nd sum} \leq 2M \cdot P\left(\left|\frac{X}{n} - t\right| \geq \delta\right)$$

[\because for Bin. r.v.
 $E(X) = nt$]

$$= 2M \cdot P(|X - nt| \geq n\delta)$$

$$\geq 2M \cdot P(|X - EX| \geq n\delta)$$

$$\leq 2M \cdot \frac{V(x)}{n^2 \delta^2} \xrightarrow{\text{By Chebyshev's inequality}} \leq 2M \cdot \frac{n t(1-t)}{n^2 \delta^2}$$

Now $t(1-t) \leq \frac{1}{4}$

$$\leq 2M \cdot \frac{n \cdot \frac{1}{4}}{n^2 \delta^2}$$

$$= \frac{M}{2n\delta^2} < \epsilon$$

[Now take n as you wish!]

[Now we can fix δ accordingly]