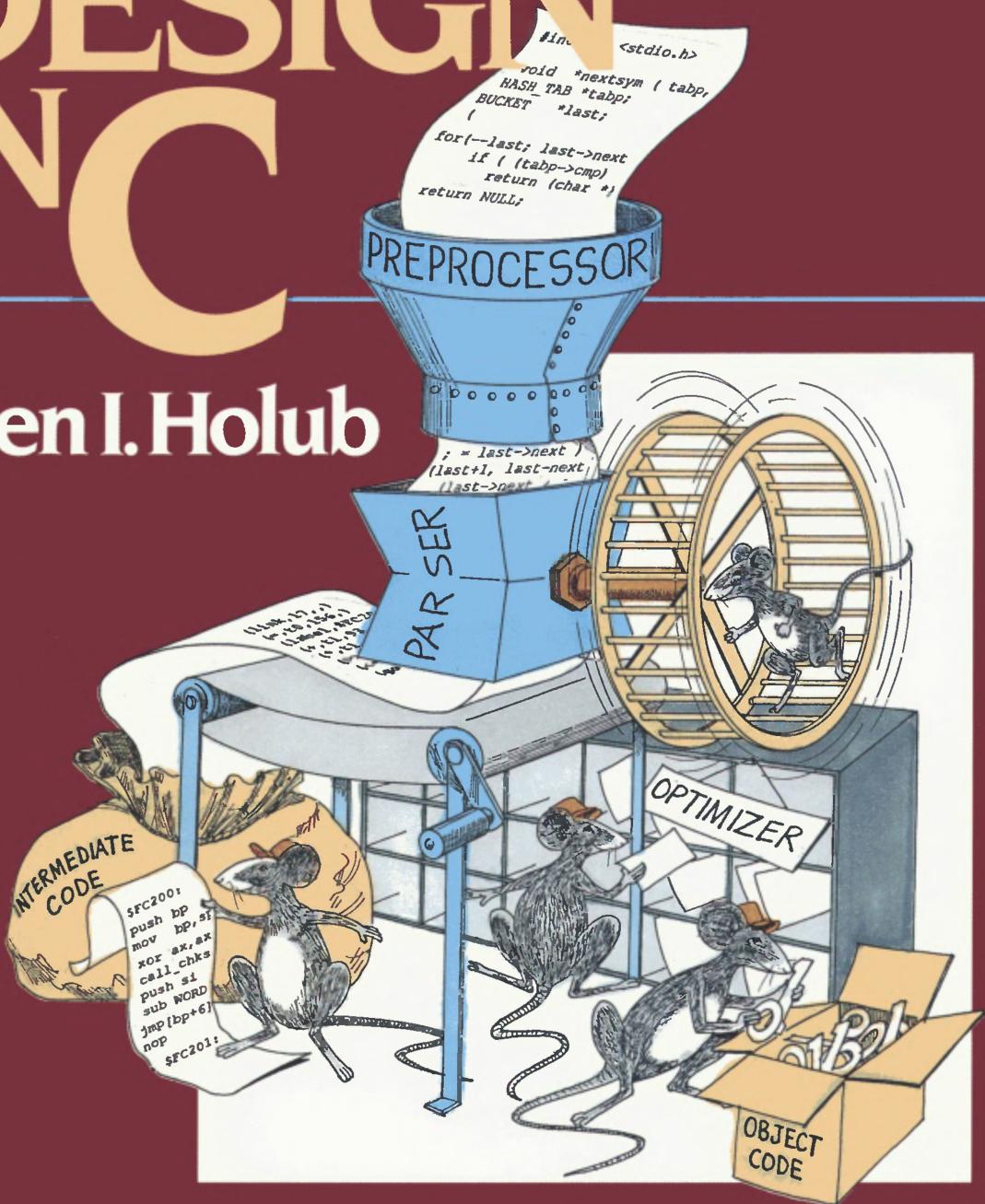


COMPILER DESIGN IN C

Allen I. Holub



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Allen I. Holub

Prentice Hall Software Series
Brian W. Kernighan, Editor



PRENTICE HALL
Englewood Cliffs, New Jersey 07632

Library of Congress Cataloging-in-Publication Data

Holub, Allen I.
Compiler design in C / Allen I. Holub.
p. cm. -- (Prentice-Hall software series)
Includes bibliographical references.
ISBN 0-13-155045-4
1. Compilers (Computer programs) 2. C (Computer program language)
I. Title. II. Series.
QA76.76.C65H65 1990
005.4'53--dc20

89-38733
CIP

Editorial/Production supervision: Kathleen Schiaparelli
Cover design: Allen I. Holub and Lundgren Graphics Ltd.
Manufacturing buyer: Margaret Rizzi

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Published by Prentice-Hall, Inc.
A division of Simon & Schuster
Englewood Cliffs, New Jersey 07632

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Printed in the United States of America

10 9 8 7 6 5 4

ISBN 0-13-155045-4

Prentice-Hall International (UK) Limited, London
Prentice-Hall of Australia Pty. Limited, Sydney
Prentice-Hall Canada Inc., Toronto
Prentice-Hall Hispanoamericana, S.A., Mexico
Prentice-Hall of India Private Limited, New Delhi
Prentice-Hall of Japan, Inc., Tokyo
Simon & Schuster Asia Pte. Ltd., Singapore
Editora Prentice-Hall do Brasil, Ltda, Rio de Janeiro

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Preface

This book presents the subject of Compiler Design in a way that's understandable to a programmer, rather than a mathematician. My basic premise is that the best way to learn how to write a compiler is to look at one in depth; the best way to understand the theory is to build tools that use that theory for practical ends. So, this book is built around working code that provides immediate practical examples of how given theories are applied. I have deliberately avoided mathematical notation, foreign to many programmers, in favor of English descriptions of the theory and using the code itself to explain a process. If a theoretical discussion isn't clear, you can look at the code that implements the theory. I make no claims that the code presented here is the only (or the best) implementation of the concepts presented. I've found, however, that looking at an implementation—at any implementation—can be a very useful adjunct to understanding the theory, and the reader is well able to adapt the concepts presented here to alternate implementations.

The disadvantage of my approach is that there is, by necessity, a tremendous amount of low-level detail in this book. It is my belief, however, that this detail is both critically important to understanding how to actually build a real compiler, and is missing from virtually every other book on the subject. Similarly, a lot of the low-level details are more related to program implementation in general than to compilers in particular. One of the secondary reasons for learning how to build a compiler, however, is to learn how to put together a large and complex program, and presenting complete programs, rather than just the directly compiler-related portions of those programs, furthers this end. I've resolved the too-many-details problem, to some extent, by isolating the theoretical materials into their own sections, all marked with asterisks in the table of contents and in the header on the top of the page. If you aren't interested in the nuts and bolts, you can just skip over the sections that discuss code.

In general, I've opted for clarity in the code rather than cleverness or theoretical efficiency. That is, since the main purpose of this book is to teach compiler-design concepts, it seemed reasonable always to go with a more understandable algorithm rather than an efficient but opaque algorithm. For example, I've used Thompson's Construction to make DFA's in LEX rather than the more direct approach recommended in Aho's book, because Thompson's construction is more understandable (it also introduces certain key concepts [like closure], in a relatively easy-to-understand context). My method for

computing LALR(1) lookaheads is also less efficient than it could be, but the algorithm is understandable and fast enough. I usually point the reader to the source for the more efficient algorithms, should he or she want to implement them.

In a sense, this book is really an in-depth presentation of several, very well documented programs: the complete sources for three compiler-generation tools are presented, as is a complete C compiler. (A lexical-analyzer generator modeled after the UNIX **lex** utility is presented along with two **yacc**-like compiler compilers.) As such, it is more of a compiler-engineering book than are most texts—a strong emphasis is placed on teaching you how to write a real compiler. On the other hand, a lot of theory is covered on the way to understanding the practice, and this theory is central to the discussion. Though I've presented complete implementations of the programs as an aid to understanding, the implementation details aren't nearly as important as the processes that are used by those programs to do what they do. It's important that you be able to apply these processes to your own programs.

The utilities are designed to be used as a learning aid. For example, **LLama** and **occs** (the two compiler compilers) can create an interactive window-oriented debugging environment, incorporating a “visible parser” that lets you watch the parse process in action. (One window shows the state and value stacks, others show the input, output, and a running commentary of what the parser's doing.) You can see the parse stack grow and shrink, watch attributes being inherited and synthesized, set breakpoints on certain conditions (a token being input, reduction by a particular production, a certain symbol on the top of stack, and so forth), and if necessary, log the entire parse (including snapshots of the stack) to a file. I've found that actually watching a bottom-up parse in action helps considerably in understanding how the parse process works, and I regularly use this visible parser in the classroom to good effect.

The C Compiler presented in Chapter Six implements an ANSI-compatible subset of C—I've left out a few things like floating point which would make the chapter even larger than it is without adding any significant value. I have tried to cover all of the hard implementation details that are usually omitted from books such as the present one. For example, the complete declaration syntax of C is supported, including structures and declarations of arbitrary complexity. Similarly, block nesting of declarations and the more complex control structures, such as switches, are also covered.

All the source code presented here is ANSI C, and is all portable to UNIX as well. For example, window management is done on the IBM-PC by emulating a standard UNIX window-management package (*curses*). The complete source code for the emulator is provided. All of the software is available electronically—versions are available for the UNIX, MS-DOS, and Macintosh environments.

I'm assuming throughout this book that you will actually read the code as well as the prose descriptions of the code. I don't waste space in the text discussing implementation details that are described adequately by the code itself. I do, however, use a lot of space describing the nonobvious parts of the programs.

The primary prerequisite for using this book is a thorough knowledge of ANSI C in particular and programming in general. You should be familiar both with the language itself and with the standard library functions described in the ANSI standard.

I've used structured programming techniques throughout the book, and have made heavy use of data abstraction and similar techniques, but I haven't described what these techniques are or why I'm using them. The more complicated data structures are explained thoroughly in the text, but a previous knowledge of basic data structures like stacks and binary trees is assumed throughout. Similarly, a knowledge of basic set theory and a familiarity with graphs is also useful. Finally, familiarity with assembly-language concepts (like how a subroutine-call works) is mandatory, but an in-depth knowledge of a specific assembly language is not required (because I've used a C subset

Prerequisites

for generated code rather than assembly language).

Though a knowledge of C is mandatory, a knowledge of UNIX or MS-DOS is not. This book is UNIX oriented only in that several UNIX tools are constructed. The tools were all developed on an IBM-PC and run nicely in that environment. By the same token, several MS-DOS implementation details and portability concerns are discussed in depth here, but I've been careful to make the code itself as portable as possible. The only potential confusion for non-UNIX users is in Appendixes D and E, where differences between my own programs and the UNIX versions are occasionally mentioned in footnotes. Just ignore these notes if you're not going to use the UNIX tools.

The book is organized so that you can use it in two ways. If you have no interest in theory and just want to build a compiler, an overview of compiler design in general is presented in Chapter One, instructions for using the compiler construction tools (`LEX` and `OCES`) are in Appendixes D and E, and code generation is discussed in Chapter Six. With these chapters behind you, you can get to the business of writing a compiler immediately, and go back later and absorb the theory.

A more rigorous approach goes through the book sequentially. You don't have to read every word of every chapter—if you're not interested in the nuts-and-bolts, just skip past those sections that describe the programs' inner workings. I'd strongly suggest reading the code, however, and reading Appendixes D and E before leaping into the text that describes these programs.

Various support functions that are used internally by the programs are concentrated in Appendix A. Covered functions do set manipulation, window management, and so forth. I've put them in one place so that the rest of the book isn't cluttered with function descriptions that aren't germane to the discussion at hand. This appendix shouldn't be viewed as an adjunct, however. The functions described in it are used heavily throughout the rest of the book, and you should be familiar with them—or at least with their calling syntax—before trying to read any of the other code.

One major organizational issue is the positioning of the theoretical and practical parts of the book. It's tempting to put all the theoretical material together at the head of each chapter so that you can skip past implementation details if you're not interested. I've opted, however, to intermix theoretical material with actual code because an examination of the code can often clarify things on the theoretical side. As I mentioned earlier, I've resolved the problem, somewhat, by isolating the theoretical material into individual sections. I often discuss the theory in one section and put a practical implementation in the following section. The theoretical sections in chapters that mix theoretical and practical material are marked with asterisks in the table of contents. This way you can skip past the implementation-related material with little difficulty, should you desire to do so.

A few compiler-related subjects are not covered here. Optimization is not discussed beyond the overview presented in Chapter Seven—I discuss how various optimizations move the code around, but I don't discuss the mechanics of optimization itself beyond a few, simple examples. Similarly, only the most common parse strategies are discussed (operator-precedence parsing is not mentioned, for example). All the material usually covered in an upper-division, undergraduate compiler-design course is covered here.

All the code in this book is written in ANSI C (I've used the Microsoft C compiler, version 5.1, for development purposes). For the most part, the MS-DOS code can be converted to the UNIX compiler by changing a few `#defines` before compiling. The disk that contains the code (see below) is shipped with the source of a small UNIX preprocessor that handles other conversion details (it does string concatenation, token pasting, etc.), and the output of this preprocessor can be submitted directly to UNIX `cc`. I'm assuming that the UNIX compiler supports those ANSI features that are implemented in most UNIX systems (like structure assignment).

Organization

Source Code and Portability

If you intend to use the code directly (without UNIX preprocessing), you'll need an ANSI-compatible compiler that supports function prototypes, and so forth. In particular:

- `<stdarg.h>` is used for variable-argument lists.
- white space around the `#` in a preprocessor directive must be permitted.
- structure assignment is used.
- `unsigned char` must be supported.
- function prototypes are used heavily.
- `isdigit()`, etc., may not have side effects.
- string concatenation is used in a few places.
- 16-character names must be permitted.

My only deviation from strict ANSI is in name lengths. In theory ANSI allows only six characters in an external name but I'm assuming that 16 are supported. I am also using the old, Kernighan & Ritchie style, subroutine-declaration syntax:

```
lorenzo( arg1, arg2 )
char *arg1;
double *arg2;
```

rather than:

```
lorenzo( char *arg1, double *arg2 )
```

I've done this because many compilers do not yet support the new syntax, and the old syntax is still legal in the standard (even though it's declared obsolescent).

I've deliberately avoided using special features of the Microsoft compiler: I've ignored things like the existence of the `far` keyword and `huge` pointers in favor of using the compact model even though the foregoing would make some of the code more efficient. By the same token, I haven't used `register` variables because the Microsoft compiler does a better job of assigning registers than I can do myself.

Unfortunately, the 8086 has an architecture that forces you to worry about the underlying machine on a regular basis, so the code has a certain amount of 8086-specific details. All of these details are isolated into macros, however, so that it's easy to port the code to a different environment by changing the macros and recompiling. I do discuss the foibles of the 8086 here and there; just skip over this material if you're not interested.

All of the source code in this book—along with executable versions of LEX, LLama, and occs—is available on disk from:

Software Engineering Consultants
 P.O. Box 5679
 Berkeley, California 94705
 (415) 540-7954

The software is available right now for the IBM-PC and UNIX. (The UNIX version is shipped on an IBM-PC, 5-1/4 disk, however. You'll have to upload it using KERMIT or some other file-transfer protocol. It has been tested under UNIX System V, BSD 4.3—I can't vouch for any other UNIX variant.) The cost is \$60.00 by a check or money order drawn on a U.S. bank. Please specify the disk size (5 $\frac{1}{4}$ " or 3 $\frac{1}{2}$ "). California residents must add local sales tax. **No purchase orders or credit cards** (sorry). A Macintosh version will be available eventually. Binary site licenses are available for educational institutions.

Bug Reports and Electronic Mail

The code in this book is bound to have a few bugs in it, though I've done my best to test it as thoroughly as possible. The version distributed on disk will always be the most recent. If you find a bug, please report it to me, either at the above address or electronically. My internet address is `holub@violet.berkeley.edu` CompuServe users can access internet from the email system by prefixing this address with `>INTERNET:` —type `help internet` for information. My UUCP address is `...!ucbvax!violet!holub`.

The UNIX USENET network is the official channel for bug fixes and general discussion of the material in this book. The *comp.compilers* newsgroup should be used for this purpose. USENET messages have a way of filtering over to other networks, like BIX, but the best way to get up-to-date information is via USENET itself. Most universities are connected to this network, and you can get access to it by getting an account on a machine at a local university. (Most schools have a mechanism for people in the community to get such accounts.) I'd prefer for all postings to be sent to me—I'll digest them and post them to *comp.compilers* via its moderator. If you want to make a submission to *comp.compilers* directly, you have to mail it to the moderator, who will post it to the network. Type *help bboard usenet moderators* to get his or her name.

This book was written largely because my students found the “standard” text—Alfred Aho, Ravi Sethi, and Jeffrey Ullman’s excellent, but at times abstruse *Compilers: Principles, Techniques, and Tools*—to be too theoretically oriented. The current volume owes a lot to Aho et al, however. I’ve used many of their algorithms, and their insights into the compiler-design practice are invaluable. I’m also indebted to Mike Lesk, Eric Schmidt, and Steve Johnson, the creators of UNIX’s **lex** and **yacc** utilities, after which the programs in this book are modeled. My neighbor, Bill Wong, provided invaluable comments on the early drafts of this book, as did many of my students. Finally, I’m grateful to Brian Kernighan, Johnson M. Hart, Andrew Appel, Norman C. Hutchinson, and N.H. Madhavji (of Bell Labs, Boston University, Princeton University, The University of Arizona, and McGill University respectively) all of whom reviewed this book before it went to press. Their comments and suggestions have made this a much better book. I am particularly indebted to Brian Kernighan, whose careful scrutiny of the entire book—both the text and the code—caught many errors that otherwise would have made it into print.

Acknowledgments

Allen Holub
Berkeley, California

This book was typeset on an IBM PC/AT using EROFF™, a version of the UNIX **troff** typesetter ported to MS-DOS by the Elan Computer Group. PostScript Times Roman and Italic were used for the text, Helvetica for chapter headings, and Courier, Courier Bold, and Courier Italic for the listings. Page proofs were generated using an Apple LaserWriter, and the final typesetting was done on a Linotronic phototypesetter using EROFF-generated PostScript. The following command line was used throughout:

```
arachne file... | autopic | tbl | troff -mm
```

The arachne preprocessor is a version of Knuth’s WEB documentation system that’s tailored for C and **troff** (rather than Pascal and **T_EX**). It runs under MS-DOS on an IBM-PC. With it, you can put the code and documentation together in a single input file. Used one way, it extracts the code and writes it out to the correct files for compilation. In a second mode it processes the code for **troff**, performing the necessary font changes, and so forth, needed to “pretty print” the code. It also automatically generates index entries for subroutine declarations. It adds line numbers to the listings and lets you reference

Typesetting Notes

these line numbers symbolically from the text (that is, you can add lines to the listings and the line numbers in the text automatically adjust themselves). Finally, it lets you discuss global variables and so forth where they're used, because it automatically moves them to the top of the output C program.

The second preprocessor, **autopic**, translates drawings generated by two commercially available drafting programs (AutoCad™ and AutoSketch™) into **troff** graphics primitives. It is much more useful than **pic** in that you have both a WYSIWYG capability and a much more powerful drawing system at your disposal. Since **troff** commands are generated as **autopic** output, the drawings are readily portable to any **troff** system.

Autopic and **arachne** are both compilers, and as such serve as an example of how you can apply the techniques presented in this book to applications other than writing compilers for standard programming languages. MS-DOS versions of **autopic** and **arachne** are available from Software Engineering at the address given earlier. Write for details.

Basic Concepts

This chapter introduces the basic concepts of compiler design. I'll discuss the internal organization of a compiler, introduce formal grammars and parse trees, and build a small recursive-descent expression compiler. Before leaping into the text, however, a word of encouragement, both about this chapter and the book in general, seems in order. Compilers are not particularly difficult programs to understand once you're familiar with the structure of a compiler in a general sort of way. The main problem is not that any one part of a compiler is hard to understand; but, rather, that there are so many parts—and you need to have absorbed most of these parts before any of them make sense. For now, my advice is to forge ahead without trying to figure out how it all ties together. You'll find that you will eventually reach a “click point” where the system as a whole suddenly makes sense.

1.1 The Parts of a Compiler

Compilers are complex programs. As a consequence, they're often broken into several distinct chunks, called *passes*, that communicate with one another via temporary files. The passes themselves are only part of the compilation process, however. The process of creating an executable image from a source-code file can involve several stages other than compilation (preprocessing, assembly, linking, and so forth). In fact, some operating systems (such as Microsoft's OS/2) can delay the final creation of an executable image until a program is actually loaded at run-time. The situation is muddled further by driver programs like UNIX's cc or Microsoft C's cl, which hide a good deal of the compilation process from you. These driver programs act as executives, controlling the various component programs that make up the compiler in such a way that you don't know that the components are being used. For the purposes of this book, I'll define a compiler as a program or group of programs that translates one language into another—in this case the source code of a high-level computer language is translated into assembly language. The assembler, linker, and so forth are not considered to be part of the compiler.

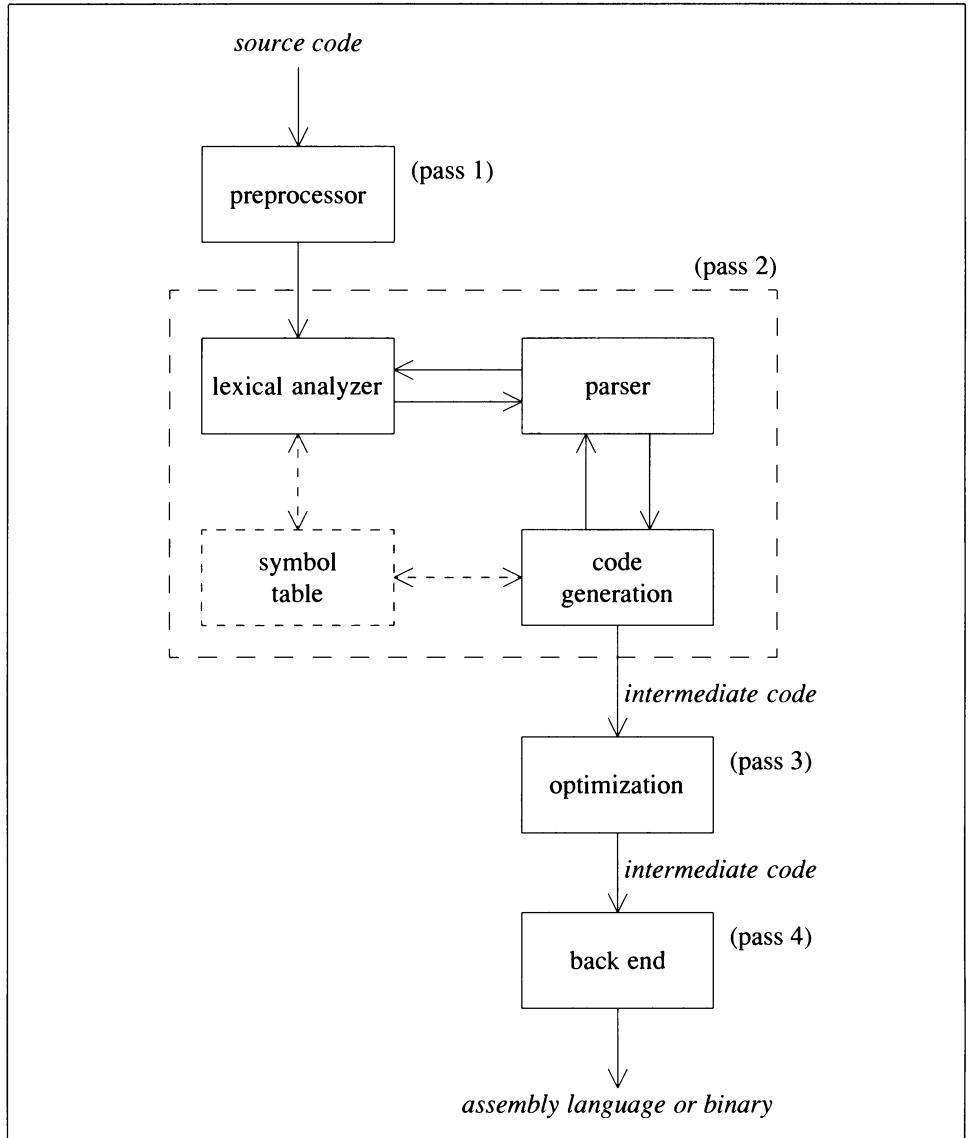
The structure of a typical four-pass compiler is shown in Figure 1.1. The preprocessor is the first pass. Preprocessors typically do macro substitutions, strip comments from the source code, and handle various housekeeping tasks with which you don't want to

Compiler passes.

Compiler defined.

Structure of a four-pass compiler.

Figure 1.1. Structure of a Typical Four-Pass Compiler



The back end.

burden the compiler proper. The second pass is the heart of the compiler. It is made up of a **lexical analyzer, parser, and code generator**, and it translates the source code into an intermediate language that is much like assembly language. The third pass is the optimizer, which improves the quality of the generated intermediate code, and the fourth pass, the **back end**, translates the optimized code to real assembly language or some form of binary, executable code. Of course, there are many variations to this structure. Many compilers don't have preprocessors; others generate assembly language in the second pass, optimize the assembly language directly, and don't have a fourth pass; still others generate binary instructions directly, without going through an ASCII intermediate language like assembler.

This book concentrates on the second pass of our model. There are several operations here too, but they interact in more complicated ways than the higher-level passes, and they share data structures (such as the symbol table) as well.

1.1.1 The Lexical Analyzer

A *phase* is an independent task used in the compilation process. Typically, several phases are combined into a single pass. The *lexical analyzer* phase of a compiler (often called a *scanner* or *tokenizer*) translates the input into a form that's more useable by the rest of the compiler. The lexical analyzer looks at the input stream as a collection of basic language elements called *tokens*. That is, a token is an indivisible lexical unit. In C, keywords like `while` or `for` are tokens (you can't say `wh ile`), symbols like `>`, `>=`, `>>`, and `>>=` are tokens, names and numbers are tokens, and so forth. The original string that comprises the token is called a *lexeme*. Note that there is not a one-to-one relationship between lexemes and tokens. A **name** or **number** token, for example, can have many possible lexemes associated with it; a `while` token always matches a single lexeme. The situation is complicated by tokens that overlap (such as the `>`, `>=`, `>>`, and `>>=`, used earlier). In general, a lexical analyzer recognizes the token that matches the longest lexeme—many languages build this behavior into the language specification itself. Given the input `>>`, a `shift` token is recognized rather than two greater-than tokens.

A lexical analyzer translates lexemes into tokens. The tokens are typically represented internally as unique integers or an enumerated type. Both components are always required—the token itself and the lexeme, which is needed in this example to differentiate the various **name** or **number** tokens from one another.

One of the early design decisions that can affect the structure of the entire compiler is the choice of a token set. You can have a token for every input symbol, or several symbols can be merged into a single token—for example, the `>`, `>=`, `>>`, and `>>=`, can be treated either as four tokens, or as a single **comparison-operator** token—the lexeme is used to disambiguate the tokens. The former approach can sometimes make code generation easier to do. Too many tokens, however, can make the parser larger than necessary and difficult to write. There's no hard-and-fast rule as to which is better, but by the time you've worked through this book, you'll understand the design considerations and will be able to make intelligent choices. In general, arithmetic operators with the same precedence and associativity can be grouped together, type-declaration keywords (like `int` and `char`) can be combined, and so forth.

The lexical analyzer is typically a self-contained unit that interfaces with the rest of the compiler via a small number (typically one or two) of subroutines and global variables. The parser calls the lexical-analyzer every time it needs a new token, and the analyzer returns that token and the associated lexeme. Since the actual input mechanism is hidden from the parser, you can modify or replace the lexical analyzer without affecting the rest of the compiler.

1.1.2 The Parser

Compilers are language translators—they translate a high-level language like C into a low-level language like 8086 assembler. Consequently, a good deal of the theoretical side of the subject is borrowed from linguistics. One such concept is the idea of parsing. To *parse* an English sentence is to break it up into its component parts in order to analyze it grammatically. For example, a sentence like this:

Jane sees Spot run.

is broken up into a subject (“Jane”) and a predicate (“sees Spot run”). The predicate is in turn broken up into a verb (“sees”), a direct object (“Spot”), and a participle that modifies the direct object (“run”). Figure 1.2 shows how the sentence is represented by a conventional sentence diagram like the ones you learned to make in the sixth grade.

A compiler performs this same process (of decomposing a sentence into its component parts) in the *parser* phase, though it usually represents the parsed sentence in a

Phases.

Scanner, tokenizer.

Tokens.

Lexemes.

Lexemes are translated to tokens.

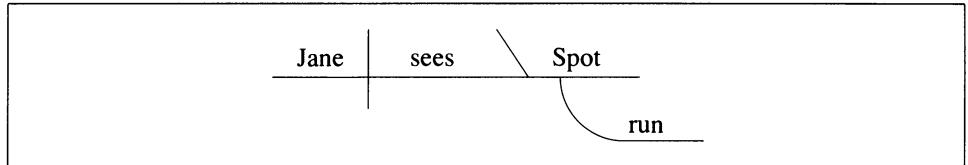
Choosing a token set.

Scanner is self-contained unit.

Parse, defined.

The parser phase.

Figure 1.2. A Sentence Diagram for *Jane Sees Spot Run*



tree form rather than as sentence diagram. (In this case, the *sentence* is an entire program.)

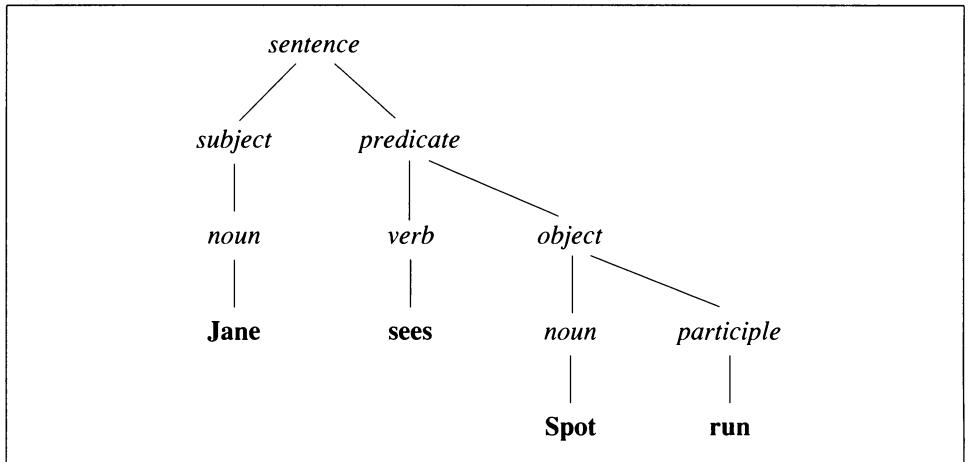
Syntax diagrams and trees.

Syntax trees.

Parse trees.

The sentence diagram itself shows the syntactic relationships between the parts of the sentence, so this kind of graph is formally called a *syntax diagram* (or, if it's in tree form, a *syntax tree*). You can expand the syntax tree, however, to show the grammatical structure as well as the syntactic structure. This second representation is called a *parse tree*. A parse tree for our earlier sentence diagram is shown in Figure 1.3. Syntax and parse trees for the expression A^*B+C^*D are shown in Figure 1.4. A tree structure is used here primarily because it's easy to represent in a computer program, unlike a sentence diagram.

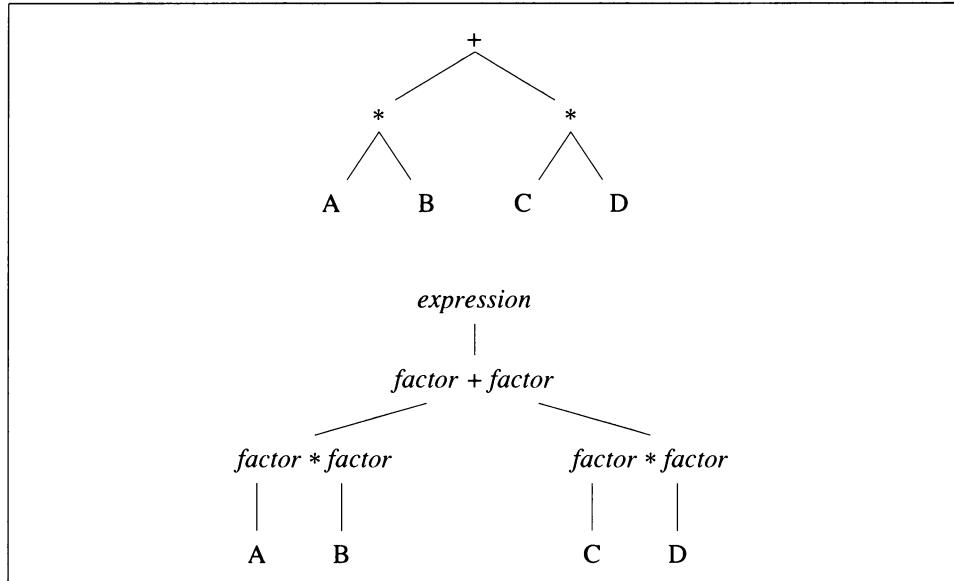
Figure 1.3. A Parse Tree for *Jane Sees Spot Run*



Sentence: formal definition.

A *sentence*, by the way, is also a technical term, though it means the same thing as it does in English. It's a collection of tokens that follow a well-defined grammatical structure. In the case of a compiler, the sentence is typically an entire computer program. The analogy is evident in a language like Pascal, which mirrors English punctuation as well as its grammar. A Pascal program ends with a period, just like an English sentence. Similarly, a semicolon is used as punctuation to separate two complete ideas, just as it separates two independent clauses in English.

To summarize: A parser is a group of subroutines that converts a token stream into a parse tree, and a parse tree is a structural representation of the sentence being parsed. Looked at another way, the parse tree represents the sentence in a hierarchical fashion, moving from a general description of the sentence (at the root of the tree) down to the specific sentence being parsed (the actual tokens) at the leaves. Some compilers create a physical parse tree, made up of structures, pointers, and so forth, but most represent the

Figure 1.4. Syntax and Parse Trees for $A^*B + C^*D$ 

parse tree implicitly.¹ Other parse methods just keep track of where they are in the tree, without creating a physical tree (we'll see how this works shortly). The parse tree itself is a very useful concept, however, for understanding how the parse process works.

1.1.3 The Code Generator

The last part of the compiler proper is the code generator. It's somewhat misleading to represent this phase as a separate component from the parser proper, because most compilers generate code as the parse progresses. That is, the code is generated by the same subroutines that are parsing the input stream. It is possible, however, for the parser to create a parse tree for the entire input file, which is then traversed by a distinct code generator, and some compilers work in this way. A third possibility is for the parser to create an intermediate-language representation of the input from which a syntax tree can be reconstructed by an optimization pass. Some optimizations are easier to perform on a syntax tree than on a linear instruction stream. A final, code-generation pass can traverse the optimizer-modified syntax tree to generate code.

Though compilers can generate object code directly, they often defer code generation to a second program. Instead of generating machine language directly, they create a program in an *intermediate language* that is translated by the compiler's *back end* into actual machine language. You can look at an intermediate language as a sort-of super assembly language that's designed for performing specific tasks (such as optimization). As you might expect, there are many flavors of intermediate languages, each useful in different applications.

There are advantages and disadvantages to an intermediate-language approach to compiler writing. The main disadvantage is lack of speed. A parser that goes straight from tokens to binary object code will be very fast, since an extra stage to process the intermediate code can often double the compile time. The advantages, however, are

Intermediate languages,
back ends.

Advantages and disad-
vantages of intermediate
languages.

1. A physical parse tree is useful for some kinds of optimizations, discussed further in Chapter Seven.

usually enough to justify the loss of speed. These are, in a nutshell, optimization and flexibility. A few optimizations, such as simple constant folding—the evaluation of constant expressions at compile time rather than run time—can be done in the parser. Most optimizations, however, are difficult, if not impossible, for a parser to perform. Consequently, parsers for optimizing compilers output an intermediate language that's easy for a second pass to optimize.

Intermediate languages give you flexibility as well. A single lexical-analyzer/parser front end can be used to generate code for several different machines by providing separate back ends that translate a common intermediate language to a machine-specific assembly language. Conversely, you can write several front ends that parse several different high-level languages, but which all output the same intermediate language. This way, compilers for several languages can share a single optimizer and back end.

Interpreters.

A final use of an intermediate language is found in incremental compilers or interpreters. These programs shorten the development cycle by executing intermediate code directly, rather than translating it to binary first, thereby saving the time necessary for assembling and linking a real program. An interpreter can also give you an improved debugging environment because it can check for things like out-of-bounds array indexing at run time.

C Code: The intermediate language used in this book.

The compiler developed in Chapter Six uses an intermediate language for the output code. The language itself is described in depth in that chapter, but some mention of it is necessary here, because I'll be using the language informally for code-generation examples throughout this book. Put simply, the intermediate language is a subset of C in which most instructions translate directly to a small number of assembly-language instructions on a typical machine (usually one or two). For example, an expression like $x=a+b*c+d$ is translated into something like this:

```
t0 = _a ;
t1 = _b ;
t1 *= _c ;
t0 += t1 ;
t0 += _d ;
```

Anonymous temporaries.

The `t0` and `t1` in the foregoing code are temporary variables that the compiler allocates to hold the result of the partially-evaluated expression. These are called *anonymous temporaries* (often shortened to just *temporaries*) and are discussed in greater depth below. The underscores are added to the names of declared variables by the compiler so that they won't be confused with variables generated by the compiler itself, such as `t0`, and `t1` (which don't have underscores in their names).

Since the intermediate language is so C like, I'm going to just use it for now without a formal language definition. Remember, though, that the intermediate language is not C (there would be little point in a compiler that translated good C into bad C)—it is really an assembly language with a C-like syntax.

1.2 Representing Computer Languages

A compiler is like every other program in that some sort of design abstraction is useful when constructing the code. Flow charts, Warnier-Orr diagrams, and structure charts are examples of a design abstraction. In compiler applications, the best abstraction is one that describes the language being compiled in a way that reflects the internal structure of the compiler itself.

1.2.1 Grammars and Parse Trees

The most common method used to describe a programming language in a formal way is also borrowed from linguistics. This method is a formal grammar, originally developed by M.I.T.'s Noam Chomsky and applied to computer programs by J.W. Backus for the first FORTRAN compilers.

Formal grammars are most often represented in a modified *Backus-Naur Form* (also called *Backus-Normal Form*), BNF for short. A strict BNF representation starts with a set of tokens, called *terminal symbols*, and a set of definitions, called *nonterminal symbols*. The definitions create a system in which every legal structure in the language can be represented. One operator is supported, the ::= operator, translated by the phrase “is defined as” or “goes to.” For example, the following BNF rule might start a grammar for an English sentence:

sentence ::= *subject predicate*

A *sentence* is defined as a *subject* followed by a *predicate*. You can also say “a sentence goes to a subject followed by a predicate.” Each rule of this type is called a *production*. The nonterminal to the left of the ::= is the *left-hand side* and everything to the right of the ::= is the *right-hand side* of the production. In the grammars used in this book, the left-hand side of a production always consists of a single, nonterminal symbol, and every nonterminal that’s used on a right-hand side must also appear on a left-hand side. All symbols that don’t appear on a left-hand side, such as the tokens in the input language, are *terminal symbols*.

A real grammar continues with further definitions until all the terminal symbols are accounted for. For example, the grammar could continue with:

subject ::= *noun*
noun ::= **JANE**

where **JANE** is a terminal symbol (a token that matches the string “Jane” in the input).

The strict BNF is usually modified to make a grammar easier to type, and I'll use a modified BNF in this book. The first modification is the addition of an OR operator, represented by a vertical bar (|). For example,

noun ::= **JANE**
noun ::= **DICK**
noun ::= **SPOT**

is represented as follows:

noun ::= **DICK | JANE | SPOT**

Similarly, a → is often substituted for the ::= as in:

Modified BNF: the | operator.

noun → **DICK | JANE**

I use the → in most of this book. I also consistently use italics for nonterminal symbols and boldface for terminals (symbols such as + and * are also always terminals—they'll also be in boldface but sometimes it's hard to tell.)

There's one other important concept. Grammars must be as flexible as possible, and one of the ways to get that flexibility is to make the application of certain rules optional. A rule like this:

article → **THE**

says that **THE** is an *article*, and you can use that production like this:

→ used for ::=

Terminals=**boldface**.
 nonterminals=*italic*.

Optional rules and ε.

Epsilon (ϵ) productions.

object → *article noun*

In English, an *object* is an *article* followed by a *noun*. A rule like the foregoing requires that all nouns that comprise an object be preceded by a participle. But what if you want the *article* to be optional? You can do this by saying that an *article* can either be the noun “the” or an empty string. The following is used to do this:

article → THE | ϵ

The ϵ (pronounced “epsilon”) represents an empty string. If the **THE** token is present in the input, then the

article → **THE**

production is used. If it is not there, however, then the *article* matches an empty string, and

article → ϵ

is used. So, the parser determines which of the two productions to apply by examining the next input symbol.

A grammar that recognizes a limited set of English sentences is shown below:

<i>sentence</i>	→	<i>subject predicate</i>
<i>subject</i>	→	<i>noun</i>
<i>predicate</i>	→	<i>verb object</i>
<i>object</i>	→	<i>noun opt_participle</i>
<i>opt_participle</i>	→	<i>participle</i> ϵ
<i>noun</i>	→	SPOT JANE DICK
<i>participle</i>	→	RUN
<i>verb</i>	→	SEES

Recognizing a sentence using the grammar.

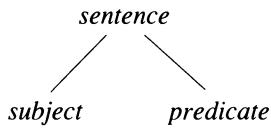
An input sentence can be recognized using this grammar, with a series of replacements, as follows:

- (1) Start out with the topmost symbol in the grammar, the goal symbol.
- (2) Replace that symbol with one of its right-hand sides.
- (3) Continue replacing nonterminals, always replacing the leftmost nonterminal with its right-hand side, until there are no more nonterminals to replace.

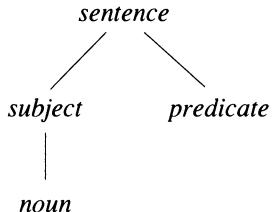
For example, the grammar can be used to recognize “Jane sees Spot run” as follows:

<i>sentence</i>	apply <i>sentence</i> → <i>subject predicate</i> to get:
<i>subject predicate</i>	apply <i>subject</i> → <i>noun</i> to get:
<i>noun predicate</i>	apply <i>noun</i> → JANE to get:
JANE predicate	apply <i>predicate</i> → <i>verb object</i> to get:
JANE verb object	apply <i>verb</i> → SEES to get:
JANE SEES object	apply <i>object</i> → <i>noun opt_participle</i> to get:
JANE SEES noun opt_participle	apply <i>noun</i> → SPOT to get:
JANE SEES SPOT opt_participle	apply <i>opt_participle</i> → <i>participle</i> to get:
JANE SEES SPOT participle	apply <i>participle</i> → RUN to get:
JANE SEES SPOT RUN	done—there are no more nonterminals to replace

These replacements can be used to build the parse tree. For example, replacing *sentence* with *subject predicate* is represented in tree form like this:



The second replacement, of *subject* with *noun*, would modify the tree like this:



and so forth. The evolution of the entire parse tree is pictured in Figure 1.5.

A glance at the parse tree tells you where the terms terminal and nonterminal come from. Terminal symbols are always leaves in the tree (they're at the end of a branch), and nonterminal symbols are always interior nodes.

Terminals are leaf nodes.
Nonterminals are interior nodes.

1.2.2 An Expression Grammar

Table 1.1 shows a grammar that recognizes a list of one or more statements, each of which is an arithmetic expression followed by a semicolon. *Statements* are made up of a series of semicolon-delimited expressions, each comprising a series of numbers separated either by asterisks (for multiplication) or plus signs (for addition).

Note that the grammar is recursive. For example, Production 2 has *statements* on both the left- and right-hand sides. There's also third-order recursion in Production 8, since it contains an *expression*, but the only way to get to it is through Production 3, which has an expression on its left-hand side. This last recursion is made clear if you make a few algebraic substitutions in the grammar. You can substitute the right-hand side of Production 6 in place of the reference to *term* in Production 4, yielding

$$\text{expression} \rightarrow \text{factor}$$

Recursion in grammar.

and then substitute the right-hand side of Production 8 in place of the *factor*:

$$\text{expression} \rightarrow (\text{expression})$$

Since the grammar itself is recursive, it stands to reason that recursion can also be used to parse the grammar—I'll show how in a moment. The recursion is also important from a structural perspective—it is the recursion that makes it possible for a finite grammar to recognize an infinite number of sentences.

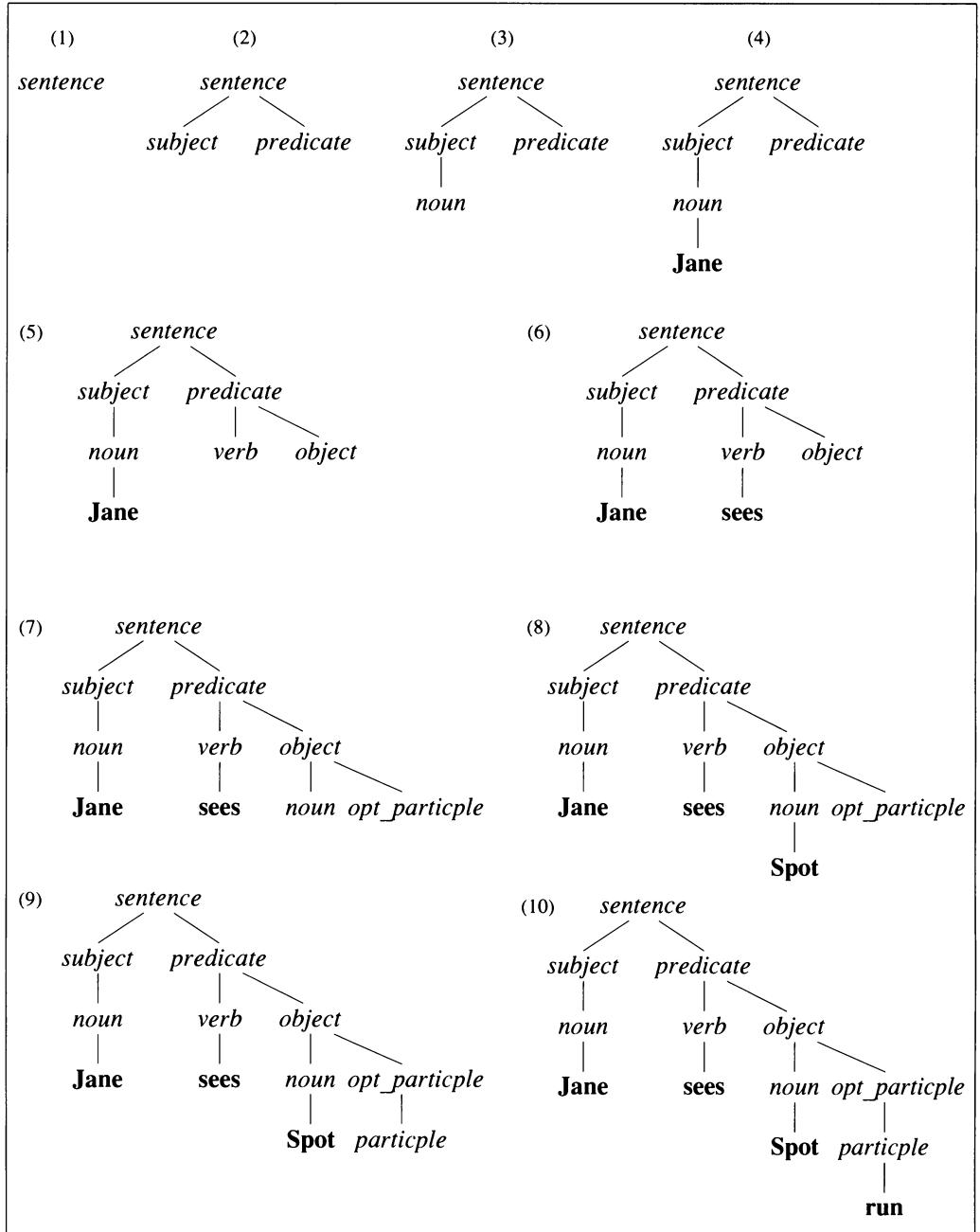
The strength of the foregoing grammar is that it is intuitive—its structure directly reflects the way that an expression goes together. It has a major problem, however. The leftmost symbol on the right-hand side of several of the productions is the same symbol that appears on the left-hand side. In Production 3, for example, *expression* appears both on the left-hand side and at the far left of the right-hand side. The property is called *left recursion*, and certain parsers (such as the recursive-descent parser that I'll discuss in a moment) can't handle left-recursive productions. They just loop forever, repetitively replacing the leftmost symbol in the right-hand side with the entire right-hand side.

Left recursion.

You can understand the problem by considering how the parser decides to apply a particular production when it is replacing a nonterminal that has more than one right-hand side. The simple case is evident in Productions 7 and 8. The parser can choose which production to apply when it's expanding a *factor* by looking at the next input symbol. If this symbol is a **number**, then the compiler applies Production 7 and replaces the

Why left recursion causes problems—an example.

Figure 1.5. Evolution of a Parse Tree



factor with a **number**. If the next input symbol was an open parenthesis, the parser would use Production 8. The choice between Productions 5 and 6 cannot be solved in this way, however. In the case of Production 6, the right-hand side of *term* starts with a *factor* which, in turn, starts with either a **number** or left parenthesis. Consequently, the parser would like to apply Production 6 when a *term* is being replaced and the next input symbol is a **number** or left parenthesis. Production 5—the other right-hand side—starts with a *term*, which can start with a *factor*, which can start with a **number** or left parenthesis, and these are the same symbols that were used to choose Production 6. To

Table 1.1. A Simple Expression Grammar

1.	<i>statements</i>	\rightarrow	<i>expression</i> ;
2.			<i>expression</i> ; <i>statements</i>
3.	<i>expression</i>	\rightarrow	<i>expression</i> + <i>term</i>
4.			<i>term</i>
5.	<i>term</i>	\rightarrow	<i>term</i> * <i>factor</i>
6.			<i>factor</i>
7.	<i>factor</i>	\rightarrow	number
8.			(<i>expression</i>)

summarize, the parser must be able to choose between one of several right-hand sides by looking at the next input symbol. It could make this decision in Productions 7 and 8, but it cannot make this decision in Productions 5 and 6, because both of the latter productions can start with the same set of terminal symbols.

The previous situation, where the parser can't decide which production to apply, is called a *conflict*, and one of the more difficult tasks of a compiler designer is creating a grammar that has no conflicts in it. The next input symbol is called the *lookahead symbol* because the parser looks ahead at it to resolve a conflict.

Unfortunately, for reasons that are discussed in Chapter Four, you can't get rid of the recursion by swapping the first and last production element, like this:

$$\textit{expression} \rightarrow \textit{term} + \textit{expression}$$

so the grammar must be modified in a very counterintuitive way in order to build a recursive-descent parser for it. Several techniques can be used to modify grammars so that a parser can handle them, and all of these are discussed in depth in Chapter Four. I'll use one of them now, however, without any real explanation of why it works. Take it on faith that the grammar in Table 1.2 recognizes the same input as the one we've been using. (I'll discuss the \vdash and ϵ that appear in the grammar momentarily.) The modified grammar is obviously an inferior grammar in terms of self-documentation—it is difficult to look at it and see the language that's represented. On the other hand, it works with a recursive-descent parser, and the previous grammar doesn't.

Table 1.2. Modified Simple-Expression Grammar

1.	<i>statements</i>	\rightarrow	\vdash
2.			<i>expression</i> ; <i>statements</i>
3.	<i>expression</i>	\rightarrow	<i>term expression'</i>
4.	<i>expression'</i>	\rightarrow	+ <i>term expression'</i>
5.			ϵ
6.	<i>term</i>	\rightarrow	<i>factor term'</i>
7.	<i>term'</i>	\rightarrow	* <i>factor term'</i>
8.			ϵ
9.	<i>factor</i>	\rightarrow	number
10.			(<i>expression</i>)

The \vdash symbol is an end-of-input marker. For the purposes of parsing, end of file is treated as an input token, and \vdash represents end of input in the grammar. In this grammar, Production 1 is expanded if the current input symbol is end of input, otherwise Production 2 is used. Note that an explicit end-of-input marker is often omitted from a

Conflicts and look-aheads.

Modified expression grammar.

End-of-input symbol (\vdash).

grammar, in which case \vdash is implied as the rightmost symbol of the starting production (the production whose left-hand side appears at the apex of the parse tree). Since eliminating the \vdash symbol removes the entire right-hand side in the current grammar, you can use the following as an alternate starting production:

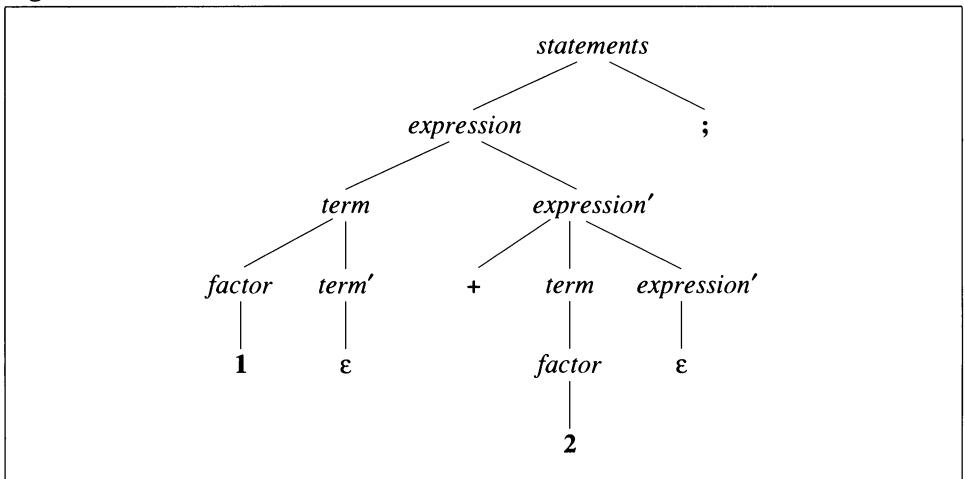
$$\text{statements} \rightarrow \epsilon \mid \text{expression} ; \text{statements}$$

In English: *statements* can go to an empty string followed by an implied end-of-input marker.

Applying ϵ .

The replacement of the left-hand side by ϵ (the empty string) occurs whenever the current input symbol doesn't match a legal lookahead symbol. In the current grammar, a *term'* is replaced with the right-hand side **factor term'* if the lookahead symbol (the next input symbol) is a *. The *term'* is replaced with ϵ if the next input symbol isn't a *. The process is demonstrated in Figure 1.6, which shows a parse of 1+2 using the modified grammar in Table 1.2. The ϵ production stops things from going on forever.

Figure 1.6. A Parse of 1+2



ϵ is a terminal, but not a token.

Note that ϵ is a terminal symbol that is not a token. It always appears at the end of a branch in the parse tree, so it is a terminal, but it does not represent a corresponding token in the input stream (just the opposite in fact—it represents the absence of a particular token in the input stream).

1.2.3 Syntax Diagrams

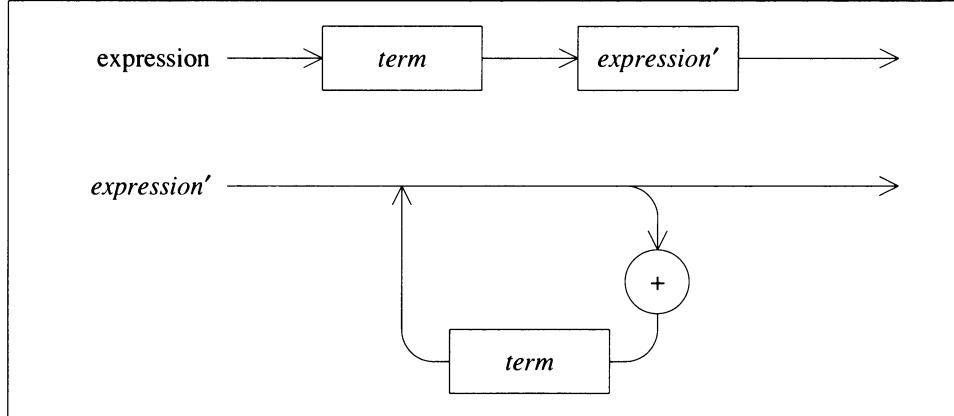
You can prove to yourself that the grammar in Table 1.2 works as expected by representing it in a different way—as a syntax diagram. We saw earlier that a syntax diagram can represent the entire syntactic structure of a parse, but you can also use it in a more limited sense to represent the syntax of a single production. Syntax diagrams are useful in writing recursive-descent compilers because they translate directly into flow charts (that's the main reason we're looking at them now). You can use them as a map that describes the structure of the parser (more on this in a moment). They are also somewhat more intuitive to an uninitiated reader, so they often make better documentation than does a formal grammar.

Translating grammars to syntax diagrams.

I'll translate our grammar into a syntax diagram in two steps. First, several of the productions can be merged together into a single diagram. Figure 1.7 represents Productions 3, 4, and 5 of the grammar in Table 1.2 on page 11. The ϵ production is represented by the uninterrupted line that doesn't go through a box. You can combine these two

graphs by substituting the bottom graph for the reference to it in the top graph, and the same process can be applied to Productions 6, 7, and 8.

Figure 1.7. Syntax Diagram for Productions 3, 4, and 5



The entire grammar in Table 1.2 is represented as a syntax diagram in Figure 1.8. The topmost diagram, for example, defines a *statement* as a list of one or more semicolon-delimited *expressions*. The same thing is accomplished by

$$\begin{array}{lcl} \textit{statements} & \rightarrow & \textit{expression} ; \\ & | & \textit{expression} ; \textit{statements} \end{array}$$

but the BNF form is harder to understand.

The merged diagram also demonstrates graphically how the modified grammar works. Just look at it like a flow chart, where each box is a subroutine, and each circle or ellipse is the symbol that must be in the input when the subroutine returns. Passing through the circled symbol removes a terminal from the input stream, and passing through a box represents a subroutine call that evaluates a nonterminal.

Diagram shows how modified grammar works.

1.3 A Recursive-Descent Expression Compiler

We now know enough to build a small compiler, using the expression grammar we've been looking at (in Table 1.2 on page 11). Our goal is to take simple arithmetic expressions as input and generate code that evaluates those expressions at run time. An expression like $a+b*c+d$ is translated to the following intermediate code:

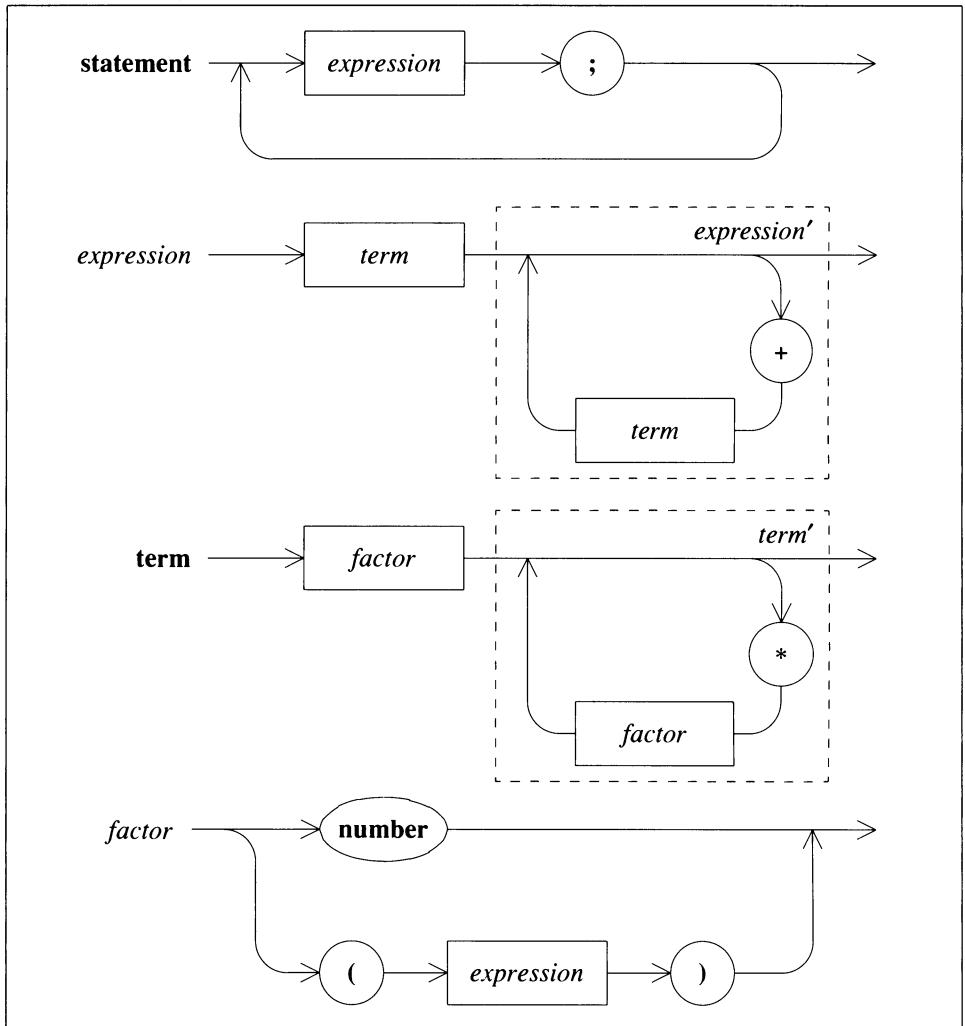
```
t0 = _a ;
t1 = _b ;
t1 *= _c ;
t0 += t1 ;
t0 += _d ;
```

1.3.1 The Lexical Analyzer

The first order of business is defining a token set. With the exception of numbers and identifiers, all the lexemes are single characters. (Remember, a *token* is an input symbol taken as a unit, a *lexeme* is the string that represents that symbol.) A **NUM_OR_ID** token is used both for numbers and identifiers; so, they are made up of a series of contiguous characters in the range '`'0'-'9'`', '`'a'-'z'`', or '`'A'-'Z'`'. The tokens themselves are defined with the macros at the top of *lex.h*, Listing 1.1. The lexical analyzer translates a semicolon into a **SEMI** token, a series of digits into a **NUM_OR_ID** token,

Expression token set.

NUM_OR_ID

Figure 1.8. A Syntax Diagram

`yytext, yylen,`

and so on. The three external variables at the bottom of *lex.h* are used by the lexical analyzer to pass information to the parser. `yytext` points at the current lexeme, which is not '\0' terminated; `yylen` is the number of characters in the lexeme; and `yylineno` is the current input line number. (I've used these somewhat strange names because both **lex** and LEX use the same names. Usually, I try to make global-variable names begin with an upper-case letter and macro names are all caps. This way you can distinguish these names from local-variable names, which are always made up of lower-case letters only. It seemed best to retain UNIX compatibility in the current situation, however.)

Simple, buffered, input system.

The lexical analyzer itself starts on line nine of *lex.c*, Listing 1.2. It uses a simple, buffered, input system, getting characters a line at a time from standard input, and then isolating tokens, one at a time, from the line. Another input line is fetched only when the entire line is exhausted. There are two main advantages to a buffered system (neither of which are really exercised here, though the situation is different in the more sophisticated input system discussed in Chapter Two). These are speed and speed. Computers like to read data in large chunks. Generally, the larger the chunk the faster the throughput. Though a 128-byte buffer isn't really a large enough chunk to make a

Listing 1.1. *lex.h*— Token Definitions and **extern** statements

```

1 #define EOI          0      /* end of input           */
2 #define SEMI         1      /* ;                      */
3 #define PLUS          2      /* +                      */
4 #define TIMES         3      /* *                      */
5 #define LP             4      /* (                      */
6 #define RP             5      /* )                      */
7 #define NUM_OR_ID     6      /* decimal number or identifier */
8
9 extern char *yytext;           /* in lex.c                */
10 extern int yyleng;
11 extern int yylineno;

```

difference, once the buffer size gets above the size of a disk cluster, the changes are more noticeable, especially if you can use the unbuffered I/O functions to do your reading and writing. The second speed issue has to do with *lookahead* and *pushback*. Lexical analyzers often have to know what the next input character is going to be without actually reading past it. They must *look ahead* by some number of characters. Similarly, they often need to read past the end of a lexeme in order to recognize it, and then *push back* the unnecessary characters onto the input stream. Consequently, there are often extra characters that must be handled specially. This special handling is both difficult and slow when you're using single-character input. Going backwards in a buffer, however, is simply a matter of moving a pointer.

Lookahead and push-back.

Listing 1.2. *lex.c*— A Simple Lexical Analyzer

```

1 #include "lex.h"
2 #include <stdio.h>
3 #include <ctype.h>
4
5 char      *yytext    = ""; /* Lexeme (not '\0' terminated)   */
6 int       yyleng     = 0;  /* Lexeme length.                 */
7 int       yylineno   = 0;  /* Input line number.            */
8
9 lex()
10 {
11     static char input_buffer[128];
12     char      *current;
13
14     current = yytext + yyleng;        /* Skip current lexeme */
15
16     while( 1 )                      /* Get the next one */
17     {
18         while( !*current )
19         {
20             /* Get new lines, skipping any leading white space on the line,
21             * until a nonblank line is found.
22             */
23
24             current = input_buffer;
25             if( !gets( input_buffer ) )
26             {
27                 *current = '\0' ;
28                 return EOI;
29             }
30

```



Listing 1.2. continued...

```

31             ++yylineno;
32
33         while( isspace(*current) )
34             ++current;
35     }
36
37     for( ; *current ; ++current )
38     {
39         /* Get the next token */
40
41         yytext = current;
42         yyleng = 1;
43
44         switch( *current )
45         {
46             case EOF: return EOI ;
47             case ';': return SEMI ;
48             case '+': return PLUS ;
49             case '*': return TIMES ;
50             case '(': return LP ;
51             case ')': return RP ;
52
53             case '\n':
54             case '\t':
55             case ' ' : break;
56
57             default:
58                 if( !isalnum(*current) )
59                     fprintf(stderr, "Ignoring illegal input <%c>\n", *current);
60                 else
61                 {
62                     while( isalnum(*current) )
63                         ++current;
64
65                     yyleng = current - yytext;
66                     return NUM_OR_ID;
67                 }
68
69             break;
70         }
71     }
72 }
73 }
```

Reading characters:
`input_buffer,`
`current.`

The input buffer used by `lex()` is declared on line 11 of Listing 1.2. `current` (on line 12) points at the current position in the buffer. On the first call, the increment on line 14 initializes `current` to point at an empty string (`yyleng` is 0 at this juncture, and `yytext` points at an empty string because of the initializer on line five). The `while` statement on line 18 tests true as a consequence. This `while` loop has two purposes. It gets lines (and increments the line number), and it skips past all blank lines (including lines that contain only white space). The loop doesn't terminate until `input_buffer` holds a nonblank line, and `current` will point at the first nonwhite character on that line.

Tokenization:

The `for` loop starting on line 37 does the actual tokenization. Single-character lexemes are recognized on lines 46–51, white space is ignored by the cases on lines 53–55, and the multiple-character `NUM_OR_ID` token is handled in the `else` clause on lines

60–67. An error message is printed if an illegal character is found. When the loop terminates, `yytext` points at the first character of the lexeme, and `yylen` holds its length.

The next time `lex()` is called, the code on line 14 adjusts the current pointer to point past the previous lexeme, and then, if the input buffer hasn't been exhausted, the `while` test on line 18 fails and you go straight to the token-isolation code. `lex()` won't get another input line until it reaches the end of the line—`*current` is '`\0`' in this case.

The remainder of `lex.c` (in Listing 1.3) addresses the problem of *lookahead*. The parser must look at the next input token without actually reading it. Though a read/pushback scheme similar to `getc()`/`ungetc()` could be used for this purpose, it's generally a good idea to avoid going backwards in the input, especially if you have to push back entire lexemes rather than single characters. The problem is solved by using two subroutines: `match(token)` evaluates to true if the next token in the input stream matches its argument—it "looks ahead" at the next input symbol without reading it. `advance()` discards the current token and advances to the next one. This strategy eliminates the necessity of a push-back subroutine such as `ungetc()`.

Solving the lookahead problem.

`match()`, `advance()`.

Lookahead.

The Lookahead variable (on line 74) holds the *lookahead* token. It's initialized to `-1`, which is not used for any of the input tokens. It's modified to hold a real token the first time `match()` is called. Thereafter, the test on line 81 will become inoperative and `match()` simply returns true if Lookahead matches its argument. This approach is relatively foolproof—though the fool in this case is myself. I know that I'll regularly forget to call an initialization routine before calling `match()`, so I'll let `match()` initialize itself the first time it's called. The `advance()` function just calls `lex()` to assign a new value to Lookahead.

Listing 1.3. `lex.c`— Match and Advance Functions

```

74 static int Lookahead = -1;           /* Lookahead token */
75
76 int match( token )
77 int token;
78 {
79     /* Return true if "token" matches the current lookahead symbol. */
80
81     if( Lookahead == -1 )
82         Lookahead = lex();
83
84     return token == Lookahead;
85 }
86
87 void advance()
88 {
89     /* Advance the lookahead to the next input symbol. */
90
91     Lookahead = lex();
92 }
```

1.3.2 The Basic Parser

Moving on to the parser, since I'm planning to start with a naive implementation and refine it, I've isolated `main()` into a small file (Listing 1.4) that I can compile once and then link with the parser proper. The parser itself is called `statements()`.

`statements()`.

The most naive parser for our grammar is shown in Listing 1.5. I've reproduced the grammar here for convenience:

Listing 1.4. main.c

```

1 main()
2 {
3     statements();
4 }
```

1.	<i>statements</i>	\rightarrow	<i>expression ;</i> \vdash
2.			<i>expression ; statement</i>
3.	<i>expression</i>	\rightarrow	<i>term expression'</i>
4.	<i>expression'</i>	\rightarrow	<i>+ term expression'</i>
5.			ϵ
6.	<i>term</i>	\rightarrow	<i>factor term'</i>
7.	<i>term'</i>	\rightarrow	<i>* factor term'</i>
8.			ϵ
9.	<i>factor</i>	\rightarrow	num_or_id
10.			(<i>expression</i>)

Subroutines correspond to left-hand sides, implement right-hand sides.

The parser generates no code, it just parses the input. Each subroutine corresponds to the left-hand side in the original grammar that bears the same name. Similarly, the structure of each subroutine exactly matches the grammar. For example, the production

$$\textit{expression} \rightarrow \textit{term expression}'$$

is implemented by the following subroutine (on line 23 of Listing 1.5):

```

expression()
{
    term();
    expr_prime();
}
```

ϵ recognized.

The ϵ production in

$$\textit{expression}' \rightarrow \textbf{PLUS} \textit{term expression}' \mid \epsilon$$

Subroutines advance past recognized tokens.

is implemented implicitly when the test on line 37 fails (if it's not a **PLUS**, it's an ϵ). Note that each subroutine is responsible for advancing past any tokens that are on the equivalent production's right-hand side.

Listing 1.5. plain.c—A Naive Recursive-Descent Expression Parser

```

1 /* Basic parser, shows the structure but there's no code generation */
2
3 #include <stdio.h>
4 #include "lex.h"
5
6 statements()
7 {
8     /* statements -> expression SEMI
9      *           | expression SEMI statements
10     */
11
12     expression();
13
14     if( match( SEMI ) )
15         advance();
```



Listing 1.5. continued ..

```
16     else
17         fprintf( stderr, "%d: Inserting missing semicolon\n", yylineno );
18
19     if( !match(EOI) )
20         statements();                                /* Do another statement. */
21     }
22
23 expression()
24 {
25     /* expression -> term expression' */
26
27     term();
28     expr_prime();
29 }
30
31 expr_prime()
32 {
33     /* expression' -> PLUS term expression'
34     *           | epsilon
35     */
36
37     if( match( PLUS ) )
38     {
39         advance();
40         term();
41         expr_prime();
42     }
43 }
44
45 term()
46 {
47     /* term -> factor term' */
48
49     factor();
50     term_prime();
51 }
52
53 term_prime()
54 {
55     /* term' -> TIMES factor term'
56     *           | epsilon
57     */
58
59     if( match( TIMES ) )
60     {
61         advance();
62         factor();
63         term_prime();
64     }
65 }
66
67 factor()
68 {
69     /* factor -> NUM_OR_ID
70     *           | LP expression RP
71     */
72 }
```



Listing 1.5. continued...

```

73     if( match(NUM_OR_ID) )
74         advance();
75
76     else if( match(LP) )
77     {
78         advance();
79         expression();
80         if( match(RP) )
81             advance();
82         else
83             fprintf( stderr, "%d: Mismatched parenthesis\n", yylineno );
84     }
85     else
86         fprintf( stderr, "%d Number or identifier expected\n", yylineno );
87 }
```

The recursion in a recursive-descent parser.

You can now see why a production like

$$\text{expression} \rightarrow \text{expression} + \text{term}$$

can't be used by a recursive-descent parser. You can implement the foregoing as follows:

```

expression()
{
    expression();

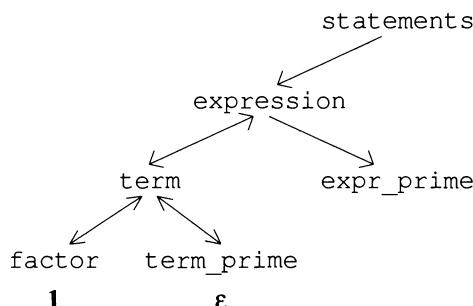
    if( !match( PLUS ) )
        error();
    else
        advance();

    term();
}
```

But the first thing that `expression()` does is call itself, the recursion never stops, and the program never terminates—at least not until it runs out of stack space and is abnormally terminated by the operating system.

At this point I'd suggest doing an exercise. Using a pencil and paper, trace what happens as the expression `1+2` is evaluated by the parser in Listing 1.5. Every time a subroutine is called, draw a downward pointing arrow and write the name of the called subroutine under the arrow; every time the subroutine returns, draw an arrow at the other end of the same line. As the parser advances past tokens, write them down under the name of the current subroutine. A partial subroutine trace for this expression is shown in Figure 1.9. The diagram shows the condition of the parser when it is in subroutine `expr_prime()` just before the `term()` call on line 40. It's advanced past the `1` and the current lookahead token is the plus sign. If you finish this diagram, an interesting fact emerges. The subroutine trace is identical to the parse tree for the same expression in Figure 1.6 on page 12. So, even though no physical parse tree is created, a parse tree is implicit in the subroutine-calling sequence.

Subroutine calling sequence mirrors the parse tree.

Figure 1.9. A Partial Subroutine Trace for 1+2

1.3.3 Improving the Parser

The naive parser discussed in the previous section is useful for explaining things, but is not much use in practice. The main difficulty is the tremendous amount of unnecessary recursion. Glancing at the syntax diagram for our grammar (in Figure 1.8 on page 14), two changes come to mind immediately. First, all the right recursion—productions in which the left-hand side also appears at the far right of the right-hand side—can be replaced by loops: If the last thing that a subroutine does is call itself, then that recursive call can be replaced by a loop. Right recursion is often called *tail* recursion.

Eliminate right recursion.

The second obvious improvement is that the same merging together of productions that was done to make the second and third graphs in Figure 1.8 can be applied to the subroutines that implement these productions. Both of these changes are made in Listing 1.6.

Merging productions into a single subroutine.

Listing 1.6. improved.c—An Improved Parser

```

1  /* Revised parser */
2
3  #include <stdio.h>
4  #include "lex.h"
5
6  void      factor      ( void );
7  void      term        ( void );
8  void      expression   ( void );
9
10 statements()
11 {
12     /* statements -> expression SEMI | expression SEMI statements */
13
14     while( !match(SEMI) )
15     {
16         expression();
17
18         if( match( SEMI ) )
19             advance();
20         else
21             fprintf( stderr, "%d: Inserting missing semicolon\n", yylineno );
22     }
23 }
24
  
```



Listing 1.6. continued...

```

25 void expression()
26 {
27     /* expression -> term expression'
28      * expression' -> PLUS term expression' | epsilon
29      */
30
31     if( !legal_lookahead( NUM_OR_ID, LP, 0 ) )
32         return;
33
34     term();
35     while( match( PLUS ) )
36     {
37         advance();
38         term();
39     }
40 }
41
42 void term()
43 {
44     if( !legal_lookahead( NUM_OR_ID, LP, 0 ) )
45         return;
46
47     factor();
48     while( match( TIMES ) )
49     {
50         advance();
51         factor();
52     }
53 }
54
55 void factor()
56 {
57     if( !legal_lookahead( NUM_OR_ID, LP, 0 ) )
58         return;
59
60     if( match(NUM_OR_ID) )
61         advance();
62
63     else if( match(LP) )
64     {
65         advance();
66         expression();
67         if( match(RP) )
68             advance();
69         else
70             fprintf( stderr, "%d: Mismatched parenthesis\n", yylineno );
71     }
72     else
73         fprintf( stderr, "%d: Number or identifier expected\n", yylineno );
74 }
```

Error recognition, FIRST sets.

I've made one additional change here as well. I've introduced a little error recovery by adding code to each subroutine that examines the current lookahead token before doing anything else. It's an error if that token cannot legitimately occur in the input. For example, expressions all have to start with either a **NUM_OR_ID** or an **LP**. If the lookahead character is a **PLUS** at the top of **expression()**, then something's wrong. This set of legitimate leading tokens is called a **FIRST** set, and every nonterminal

symbol has its own FIRST set. FIRST sets are discussed in depth in Chapter Three—the informal definition will do for now, though.

The `legal_lookahead()` subroutine in Listing 1.7 checks these FIRST sets, and if the next input symbol is not legitimate, tries to recover from the error by discarding all input symbols up to the first token that matches one of the arguments. The subroutine takes a variable number of arguments, the last one of which must be zero. I've used the ANSI variable-argument mechanism here, so the routine can take any number of arguments as parameters, but the last one must be a 0. (This mechanism, and the `<stdarg.h>` file, is described in Appendix A if you're not already familiar with it.)

One final C style note is needed to head off the inevitable criticism of the `goto` statement on line 118 of Listing 1.7. Though many programmers contend with almost religious fervor that the `goto` should be obliterated from all structured programs, I strongly feel that there are a few situations where a judiciously used `goto` makes for better code. Here, the `goto` branch to a single label is vastly preferable to multiple `return` statements. A subroutine that has a single exit point is much more maintainable than one with several exit points. You can put a single breakpoint or debugging diagnostic at the end of the subroutine instead of having to sprinkle them all over the place. You also minimize the possibility of accidentally falling off the bottom of the subroutine without returning a valid value. My rules of thumb about `gos` are as follows:

- Goto's are appropriate in two situations: (1) to eliminate multiple return statements and (2) to break out of nested loops. You can also do (2) with a flag of some sort (`while (!done)`), but flags tend to make the code both larger and harder to read, so should be avoided.
- Don't use a `goto` unless it is the only solution to the problem. You can often eliminate the need for a `goto` by reorganizing the code.
- A subroutine should have at most one label.
- All `goto` branches should be in a downwards direction to that single label.
- The target of a `goto` branch should be in the same block or at a higher (more outer) nesting level than the `goto` itself. Don't do this:

```
{
    goto label;
}
...
{
    label:
}
```

Error recovery:
`legal_lookahead()`.

Style note: the `goto`.

Listing 1.7. *improved.c*—Error Recovery for the Improved Parser

```

75 #include <stdarg.h>
76
77 #define MAXFIRST 16
78 #define SYNCH SEMI
79
80 int     legal_lookahead( first_arg )
81 int     first_arg;
82 {
83     /* Simple error detection and recovery. Arguments are a 0-terminated list of
84      * those tokens that can legitimately come next in the input. If the list is
85      * empty, the end of file must come next. Print an error message if
86      * necessary. Error recovery is performed by discarding all input symbols
87      * until one that's in the input list is found
88      *
89      * Return true if there's no error or if we recovered from the error,
90      * false if we can't recover.

```



Listing 1.7. continued...

```

91      */
92
93      va_list      args;
94      int          tok;
95      int          lookaheads[MAXFIRST], *p = lookaheads, *current;
96      int          error_printed = 0;
97      int          rval         = 0;
98
99      va_start( args, first_arg );
100
101     if( !first_arg )
102     {
103         if( match(EOI) )
104             rval = 1;
105     }
106     else
107     {
108         *p++ = first_arg;
109         while( (tok = va_arg(args, int)) && p < &lookaheads[MAXFIRST] )
110             *++p = tok;
111
112         while( !match( SYNCH ) )
113         {
114             for( current = lookaheads; current < p ; ++current )
115                 if( match( *current ) )
116                 {
117                     rval = 1;
118                     goto exit;
119                 }
120
121             if( !error_printed )
122             {
123                 fprintf( stderr, "Line %d: Syntax error\n", yylineno );
124                 error_printed = 1;
125             }
126
127             advance();
128         }
129     }
130
131     exit:
132     va_end( args )
133     return rval;
134 }
```

1.3.4 Code Generation**Recognizers.**

The parsers that we just looked at are, strictly speaking, *recognizer* programs in that, if they terminate without an error, the input sentence is a legal sentence in the grammar. All they do is recognize legal input sentences. Our goal is to build a compiler, however, and to do this, you need to add code generation to the bare-bones recognizer. Given the input

```
1 + 2 * 3 + 4
```

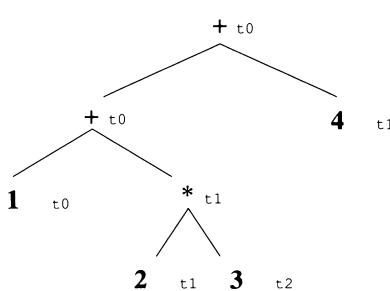
a typical compiler generates the following code:

```
t0 = 1
t1 = 2
t2 = 3
t1 *= t2
t0 += t1
t1 = 4
t0 += t1
```

An optimizer will clean up the unnecessary assignments. It's useful, for now, to look at the raw output, however. The temporary variables (t_0 , and so forth) are maintained internally by the compiler. A real compiler typically uses registers or the run-time stack for temporaries. Here, they're just global variables. The expression is evaluated operator by operator, with each temporary holding the result of evaluating the current subexpression. Sometimes (as is the case in this example) several temporaries must exist simultaneously in order to defer application of an operator because of precedence or associativity problems. Here, t_0 holds the left operand of the addition operator until the higher-precedence multiply is performed.

You can also look at the temporary-variable assignments in terms of a syntax tree. The syntax tree for $1+2*3+4$ is shown in Figure 1.10. The nodes are marked with the names of the temporaries that hold the evaluated subexpression represented by the subtree.

Figure 1.10. Syntax Tree for $1+2*3+4$



The first thing you need to generate code is a mechanism for allocating temporaries. Ideally, they should be recycled—the temporaries should be reused after they are no longer needed in the current subexpression. A simple, but effective, mechanism to do this is shown in Listing 1.8. (We'll look at a more elaborate system in Chapter Six.) A stack of temporary-variable names is declared on line one. When a new name is required, `newname()` pops one off the stack. When the temporary is no longer needed, a `freename()` call pushes it back.

The next code-generation problem is determining the name of the temporary that holds the partially-evaluated expression at any given moment. This information is passed around between the subroutines in the normal way, using arguments and return values.

To demonstrate the differences between the two methods, I'll show two parsers, one that uses return values exclusively, and another that uses arguments exclusively. The subroutine-calling tree for a parse of $1+2*3+4$; (using the improved parser) is in Figure 1.11. The subroutines that generated this tree (and the earlier code) are in Listing 1.9. The arrows are marked to show the flow of temporary-variable names during the parse. Control flow goes counterclockwise around the drawing (the nodes are subscripted to show the order in which they're visited).

Temporary variables at run time.

Temporaries on the syntax tree.

Compile-time, temporary-variable management.

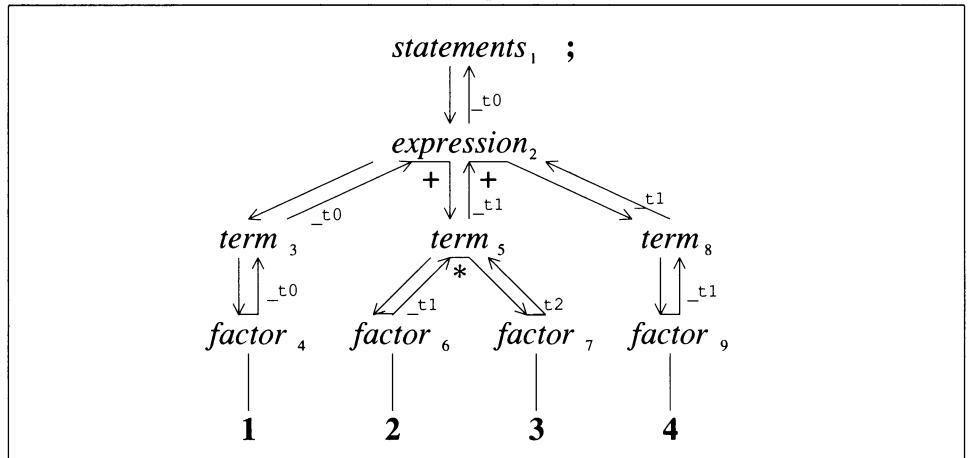
`newname()`,
`freename()`.

Using return values in parser.

Listing 1.8. *name.c*—Temporary-Variable Allocation Routines

```

1  char *Names[] = { "t0", "t1", "t2", "t3", "t4", "t5", "t6", "t7" };
2  char **Namep = Names;
3
4  char *newname()
5  {
6      if( Namep >= &Names[ sizeof(Names) / sizeof(*Names) ] )
7      {
8          fprintf( stderr, "%d: Expression too complex\n", yylineno );
9          exit( 1 );
10     }
11
12     return( *Namep++ );
13 }
14
15 freename(s)
16 char *s;
17 {
18     if( Namep > Names )
19         *--Namep = s;
20     else
21         fprintf(stderr, "%d: (Internal error) Name stack underflow\n",
22                                     yylineno );
23 }
```

Figure 1.11. A Subroutine Trace of 1+2 (Improved Parser)

Listing 1.9. *retval.c*— Code Generation Using Return Values

```
1 #include <stdio.h>
2 #include "lex.h"
3
4 char *factor ( void );
5 char *term ( void );
6 char *expression ( void );
7
8 extern char *newname( void );
9 extern void freename( char *name );
10
11 statements()
12 {
13     /* statements -> expression SEMI | expression SEMI statements */
14
15     char *tempvar;
16
17     while( !match(EOI) )
18     {
19         tempvar = expression()
20
21         if( match( SEMI ) )
22             advance();
23         else
24             fprintf( stderr, "%d: Inserting missing semicolon\n", yylineno );
25
26         freename( tempvar );
27     }
28 }
29
30 char *expression()
31 {
32     /* expression -> term expression'
33      * expression' -> PLUS term expression' | epsilon
34      */
35
36     char *tempvar, *tempvar2;
37
38     tempvar = term();
39     while( match( PLUS ) )
40     {
41         advance();
42         tempvar = term();
43         printf("    %s += %s\n", tempvar, tempvar2 );
44         freename( tempvar2 );
45     }
46
47     return tempvar;
48 }
49
50 char *term()
51 {
52     char *tempvar, *tempvar2 ;
53
54     tempvar = factor();
55     while( match( TIMES ) )
56     {
57         advance();
58         tempvar2 = factor();
59         printf("    %s *= %s\n", tempvar, tempvar2 );
```



Listing 1.9. continued...

```

60         freename( tempvar2 );
61     }
62
63     return tempvar;
64 }
65
66 char    *factor()
67 {
68     char *tempvar;
69
70     if( match(NUM_OR_ID) )
71     {
72         /* Print the assignment instruction. The %0.*s conversion is a form of
73          * %X.Ys, where X is the field width and Y is the maximum number of
74          * characters that will be printed (even if the string is longer). I'm
75          * using the %0.*s to print the string because it's not \0 terminated.
76          * The field has a default width of 0, but it will grow the size needed
77          * to print the string. The ".*" tells printf() to take the maximum-
78          * number-of-characters count from the next argument (yylen).
79         */
80
81         printf("      %s = %0.*s\n", tempvar = newname(), yylen, yytext );
82         advance();
83     }
84     else if( match(LP) )
85     {
86         advance();
87         tempvar = expression();
88         if( match(RP) )
89             advance();
90         else
91             fprintf(stderr, "%d: Mismatched parenthesis\n", yylineno );
92     }
93     else
94         fprintf( stderr, "%d: Number or identifier expected\n", yylineno );
95
96     return tempvar;
97 }
```

Generate `t0=1`.

Generate arithmetic instructions.

Using subroutine arguments to pass information.

A likely place to generate instructions of the form `t0=1` is in `factor()`, the subroutine that reads the 1. `factor()` calls `newname()` to get an anonymous-temporary name, generates the code to copy the input number into the temporary, and then returns the name to its parent. Similarly, the best place to generate multiplication instructions is the place where the times signs are read: in `term()`. After the two `factor()` calls, `tempvar` and `tempvar2` hold the names of the two temporaries. Code is generated to do the multiplication, one of the temporaries is freed, and the other (which holds the result of the multiply) is passed back up to the parent. So, this temporary, the one that holds the result of the subexpression evaluation, is used in the next step. Addition is handled the same way in `expression()`.

Just to make sure that you understand the process, I suggest taking a moment to run through a parse of `1 * (2+3) * 4` by hand, creating a subroutine-calling graph as in the previous example and watching how the code is generated.

Listing 1.10 shows the parser, modified once more to use subroutine arguments rather than return values. Here, instead of allocating the names when the instructions are generated, the temporary variables are allocated high up in the tree. Each subroutine passes

to its child the name of the temporary in which it wants the subexpression result. That is, the high level routine is saying: “Do what you must to evaluate any subexpressions, but by the time you’re finished, I want the result to be in the temporary variable whose name I’m passing to you.” Recursive-descent compilers often use both of the methods just discussed—neither is more or less desirable. The code can be easier to understand if you restrict yourself to one or the other, though that’s not always possible.

Listing 1.10. *args.c*— Code Generation Using Subroutine Arguments

```
1 #include <stdio.h>
2 #include "lex.h"
3
4 void factor ( char *tempvar );
5 void term ( char *tempvar );
6 void expression ( char *tempvar );
7
8 extern char *newname( void );
9 extern void freename( char *name );
10
11 statements()
12 {
13     /* statements -> expression SEMI | expression SEMI statements */
14
15     char *tempvar;
16
17     while( !match(EOI) )
18     {
19         expression( tempvar = newname() );
20         freename( tempvar );
21
22         if( match( SEMI ) )
23             advance();
24         else
25             fprintf( stderr, "%d: Inserting missing semicolon\n", yylineno );
26     }
27 }
28
29 void expression( tempvar )
30 char *tempvar;
31 {
32     /* expression -> term expression'
33      * expression' -> PLUS term expression' | epsilon
34      */
35
36     char *tempvar2;
37
38     term( tempvar );
39     while( match( PLUS ) )
40     {
41         advance();
42
43         term( tempvar2 = newname() );
44
45         printf("    %s += %s\n", tempvar, tempvar2 );
46         freename( tempvar2 );
47     }
48 }
```



Listing 1.10. continued...

```

50 void      term( tempvar )
51 char      *tempvar;
52 {
53     char  *tempvar2 ;
54
55     factor( tempvar );
56     while( match( TIMES ) )
57     {
58         advance();
59
60         factor( tempvar2 = newname() );
61
62         printf("    %s *= %s\n", tempvar, tempvar2 );
63         freename( tempvar2 );
64     }
65 }
66
67 void      factor( tempvar )
68 char      *tempvar;
69 {
70     if( match(NUM_OR_ID) )
71     {
72         printf("    %s = %0.*s\n", tempvar, yylen, yytext );
73         advance();
74     }
75     else if( match(LP) )
76     {
77         advance();
78         expression( tempvar );
79         if( match(RP) )
80             advance();
81         else
82             fprintf( stderr, "%d: Mismatched parenthesis\n", yylineno );
83     }
84     else
85         fprintf( stderr, "%d: Number or identifier expected\n", yylineno );
86 }
```

1.4 Exercises

- 1.1. Write a grammar that recognizes a C variable declaration made up of the following keywords:

```
int      char      long   float   double  signed  unsigned short
const    volatile
```

and a variable name.

- 1.2. Write a grammar that recognizes a C variable declaration made up only of legal combinations of the following keywords:

```
int      char      long   float   double  signed  unsigned short
const    volatile
```

and a variable name. The grammar should be able to accept all such legal declarations. For example, all the following should be accepted:

```

volatile unsigned long int x;
unsigned long volatile int x;
long unsigned volatile int x;
long volatile unsigned int x;
...

```

but something like

```
unsigned signed short long x;
```

should not be accepted. Remember that the **int** keyword is optional in a declaration.

- 1.3. Modify your solution to the previous exercise so that declarations for arrays, pointers, pointers to arrays, arrays of pointers, arrays of pointers to arrays, and so on, are also recognized. That is, all legal combination of stars, brackets, parentheses, and names should be recognized.
- 1.4. Write a grammar (and a recursive-descent compiler for that grammar) that translates an English description of a C variable into a C-style variable declaration. For example, the input:

```

x is a pointer to an array of 10 pointers to functions that return int.
y is an array of 10 floats.
z is a pointer to a struct of type a_struct.

```

should be translated to:

```

int      (*(*x)[10])();
float    y[10];
struct   a_struct *z;

```

- 1.5. Modify either of the expression compilers (in Figures 1.11 or 1.10) so that the C ++ and -- operators are supported.
- 1.6. LISP uses a prefix notation for arithmetic expressions. For example, 1+2 is represented as (+ 1 2), and 1+2*3 is (+ 1 (* 2 3)). Modify the expression compiler so that it translates infix expressions to prefix.
- 1.7. Write a LISP-to-infix translator.
- 1.8. Modify the expression compiler so that it translates expressions into postfix notation, such as that used by a Hewlett-Packard calculator. For example, the expression (1+2) * (3+4) should be translated to:

```
1 2 + 3 4 + *
```

- 1.9. Modify the expression compiler so that it prints the parse tree created by its input. I suggest creating a physical parse tree (with structures and so forth) and then printing the tree by traversing the physical parse tree.
- 1.10. (This is a very difficult problem.)
 - a. Try to write a context-free grammar that correctly parses both "time flies like an arrow" and "fruit flies like a banana."
 - b. One of the things that defines context-free grammars is that the left-hand side always consists of a single nonterminal symbol. How would the foregoing work if you were permitted to use more than one terminal or nonterminal symbol on a left-hand side? Try to write a parser for this sort of grammar.

Input and Lexical Analysis

This chapter looks at input strategies and at lexical analysis. I'll discuss a set of buffered input routines and construct L^EX, a program modeled after the UNIX `lex` utility, that translates regular expressions into a lexical analyzer. It's worth understanding how L^EX works, even if you're not going to build a version of your own. Various theoretical concepts such as finite automata and ϵ closure will crop up again when I discuss how programs like `ccs` and `yacc` generate bottom-up parse tables, and you need to be able to understand how these tables are created in order to use these programs effectively. The concepts are easier to understand in the context of L^EX, however, so the theoretical material in this chapter is actually an introduction to the concepts that you'll need later. I'm using a bootstrap approach in that L^EX itself uses a hard-coded lexical analyzer and a recursive-descent parser to do its work. As such, it's a good example of a compiler built by hand, without special tools.

The techniques used for lexical analysis are useful in many programming applications other than compilers. Efficient I/O is a concern in virtually every computer program. Similarly, lexical analyzers are pattern-recognition engines—the concepts discussed here can be applied to many programs that need to recognize patterns: editors, bibliographic data-base programs, and so forth. You can extend the techniques to do things like assembly-line quality control and network-protocol processing.

If you intend to read the implementation parts of this chapter rather than just the theory, you should read Appendix D, which contains a user's manual for L^EX, before proceeding. Also many of the support routines used by L^EX are presented in Appendix A (the set routines are used heavily in the code that follows, and the hash functions are used as well).

2.1 The Lexical Analyzer as Part of a Compiler*

The main purpose of a lexical analyzer in a compiler application is to translate the

* An asterisk appended to a section heading is used throughout this and subsequent chapters to indicate theoretical material. Implementation-oriented sections are not so marked.

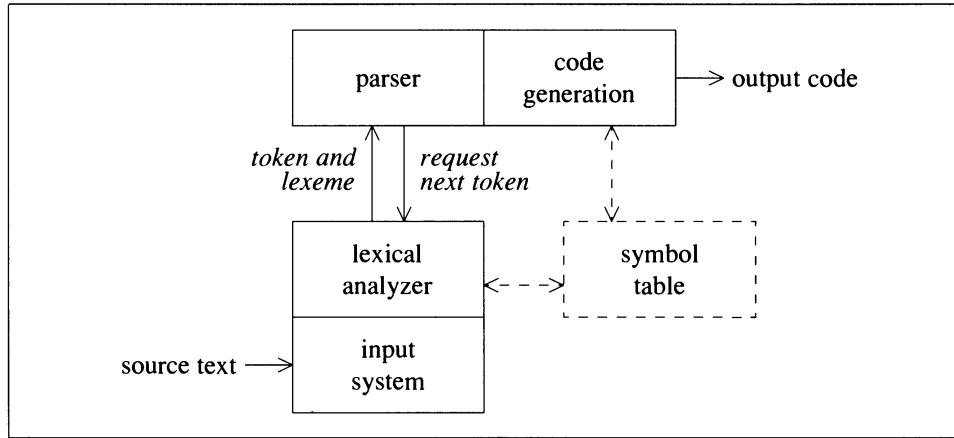
input stream into a form that is more manageable by the parser. It translates input strings or *lexemes*, into *tokens*—arbitrary integer values that represent the lexemes. A token can have a one-to-one relationship with a lexeme. For example, the keyword **while** is associated with a single token. More generic tokens such as identifiers or numbers have several lexemes associated with them. Lexical analyzers often have auxiliary functions as well. For example, a lexical analyzer can discard comments and skip over white space. Isolating this housekeeping from the parser can simplify the parser design (and the grammar of the language). The analyzer can keep track of the current line number so that intelligent error messages can be output by the parser. Program listings that show the source code intermixed with error messages are usually created by the lexical analyzer.

Lexemes, tokens.

The lexical analyzer is an independent compilation phase that communicates with the parser over a well-defined and simple interface. The relationship is pictured in Figure 2.1. The parser calls a single lexical-analyzer subroutine every time it needs a new token, and that subroutine returns the token and associated lexeme.

Interface to parser.

Figure 2.1. Interaction Between the Lexical Analyzer and Parser



This organization has several things going for it. Since it's an independent phase, the lexical analyzer is easy to maintain because changes to the analyzer do not affect the compiler as a whole, provided that the interface is not changed. Moreover, much of the code that comprises the lexical analyzer is the same for every compiler, regardless of the input language, so you can recycle much of the lexical analyzer's code. The only things that change from language to language in the table-driven lexical analyzers described later in this chapter are the tables themselves. Other advantages include speed—an independent lexical analyzer can optimize character-read times because it can read large amounts of data at once, and portability—the peculiarities of reading the source code under a particular operating system are all confined to the lexical analyzer itself. Notice that the actual input system in Figure 2.1 is isolated completely from the parser, even though it's closely linked to the lexical analyzer.

Advantages of independent lexical analyzers.

Sometimes a more complex interaction between lexical analyzer and parser is required. For example, the **typedef** statement in C effectively creates new keywords in the language. After the parser has processed the statement:

Shared symbol table.

```
typedef int alphonso;
```

the lexical analyzer must treat the input string *alphonso* as if it were a **type** token rather than as an **identifier** token. This sort of high-level communication is usually done through a shared data structure such as the symbol table. In this case, the parser can enter

`alphonso` into the symbol table, identifying it as a **typedef** name, and the lexical analyzer can check the symbol table to determine if a string is a **type** or **identifier** token.

Attributes.

A lexical analyzer can also do work other than simple pattern recognition. For example, when it reads a numeric constant, it can translate that constant into the associated number [in a manner similar to `atoi()`] and return that number along with the token and lexeme. When an identifier is read, the analyzer can look up the identifier in the symbol table and return a pointer to a symbol-table entry along with the token. These additional values associated with individual tokens are called *attributes*. (Note that the lexeme is also an attribute of the token, because it's a quantum of information that is associated with the token.) In general, it's best to restrict the lexical analyzer to simple pattern-recognition tasks in order to make it easier to maintain. If the analyzer is an independent module that performs only one task (pattern recognition), it's a simple matter to replace it if necessary.

There's one final point to make. Lexical analysis is often complicated if a language is not designed with ease-of-compilation in mind. For example,¹ PL/I keywords are not reserved—you can use them as identifiers—and the lexical analyzer has to determine what it's looking at based on surrounding context. You can say something like this in PL/I:

```
if then then then = else; else else = then;
```

Separating the keyword `then` from the identifier `then` can be quite difficult.

Discarding characters.

2.2 Error Recovery in Lexical Analysis*

It's possible, of course, for errors to occur in the lexical-analysis as well as the parsing phase. For example, the at sign (@) and backquote (`) are both illegal outside of a string in a C program. The lexical analyzer can recover from these errors in several ways, the simplest of which just discards the offending character and prints an appropriate error message. Even here, there are some choices that are driven by the application, however. If the last character of a multiple-character lexeme is incorrect, the analyzer can discard the entire malformed lexeme or it can discard only the first character of the lexeme and then try to rescan. Similarly, the lexical analyzer could try to correct the error. Some operating systems have a “do what I mean” feature that works along these lines. When faced with an error, the operating system's command-line interpreter (which is a compiler) tries to determine what the user meant to type and proceeds accordingly. If a word has only one misspelled letter, it's not too difficult to correct the problem by simple inspection.²

1. This example (and the other PL/I example, below) is borrowed from [Aho], p. 87 and p. 90.

2. There are other, more sophisticated techniques that can be used to determine the similarity of two words. The most common technique is the “soundex” algorithm developed by Margaret Odell and Robert Russel, and described in [Knuth], vol. 3, pp. 391–392. Also of interest is the Ratcliff/Obershelp algorithm, described in [Ratcliff], and Allen Bickel's algorithm, described in [Bickel] and implemented in C in [Howell].

2.3 Input Systems*

Since the input system is usually an independent module, and since the concerns here are divorced from the mechanics of recognizing tokens, I'll look at input systems in depth before moving on to the issues of lexical analysis per se.

The lowest-level layer of the lexical analyzer is the input system—the group of functions that actually read data from the operating system. For the same reason that the analyzer should be a distinct module, it's useful for the input system itself to be an independent module that communicates with the analyzer via well-defined function calls. Since the analyzer itself can be isolated from the input mechanics, the resulting code is more portable. Most of the system-dependent operations of the analyzer are concentrated into the input layer.

Issues of optimization aside, most compilers spend a good portion of their time in the lexical analysis phase. Consequently, it's worthwhile to optimize lexical analysis for speed. The standard C buffered input system is actually a poor choice for several reasons. First, most buffered systems copy the input characters at least three times before your program can use them: from the disk to a buffer maintained by the operating system, from that buffer to a second buffer that's part of the FILE structure, and finally from the FILE buffer to the string that holds the lexeme. All this copying takes both time and buffer space. Moreover, the buffer size is not optimal. The more you can read from the disk at one time, the faster your input routines tend to be (though this is operating-system dependent—there's not much advantage under UNIX in reading more than one block at a time; MS-DOS, however, performs much better with very large reads).

The other issue is lookahead and pushback. The lexical analyzer may have to look ahead several characters in the input to distinguish one token from another, and then it must push the extra characters back into the input. Consider the earlier PL/I expression:

```
if then then then = else; else else = then;
```

The lexical analyzer can distinguish the **else** keyword from the **else** identifier by looking at the characters that follow the lexeme. The **else** must be an identifier if it is followed by an equal sign, for example. Another example is a PL/I **declare** statement like this:

```
declare ( arg1, arg2, ..., argN )
```

The lexical analyzer can't distinguish the **declare** keyword from an identifier (a subroutine name in this case) until it has read past the rightmost parenthesis. It must read several characters past the end of the lexeme, and then push the extra characters back into the input stream once a decision has been made, and this lookahead and pushback must be done as efficiently as possible.

A final, admittedly contrived, example demonstrates that pushback is necessary even in recognizing individual tokens. If a language has the three tokens *xxyy*, *xx*, and *y*, and the lexical analyzer is given the input *xx* it should return an *xx* token followed by a *y* token. In order to distinguish, however, it must read at least four characters (to see if *xxyy* is present) and then push the two *y*'s back into the input stream.

Most programming languages are designed so that problems such as the foregoing are not issues. If the tokens don't overlap, you can get by with only one character of pushback. L^EX however, can't make assumptions about the structure of the lexemes, so must assume the worst.

The pushback problem means that you can't use the normal buffered input functions because `ungetc()` gives you only one character of pushback. You can add a layer around `getc()` that gives you more pushback by using a stack, as is demonstrated in Listing 2.1. Push back a character by pushing it onto the stack, and get the next input

The input system is an independent module.

Optimizing for speed.

Lookahead, pushback.

`ungetc()` inappropriate.
Stack-based pushback.

character either from the stack (if it's not empty) or from the real input stream (if it is). UNIX **lex** uses this method. A better method is described in the remainder of this section.

Listing 2.1. Using a Stack for Multiple-Character Pushback

```

1  #include <stdio.h>
2  #define SIZE 128           /* Maximum number of pushed-back characters */
3
4  /* Pbackbuf is the push-back stack.
5   * Pbackp is the stack pointer. The stack grows down, so a push is:
6   *     --Pbackp=c and a pop is: c=*Pbackp++
7   * get() evaluates to the next input character, either popping it off the
8   * stack (if it's not empty) or by calling getc().
9   * unget(c) pushes c back. It evaluates to c if successful, or to -1 if the
10  * pushback stack was full.
11 */
12
13 int Pbackbuf[SIZE];
14 int *Pbackp = &Pbackbuf[SIZE];
15
16 #define get(stream) (Pbackp < &Pbackbuf[SIZE] ? *Pbackp++ : getc(stream))
17 #define unget(c)    (Pbackp <= Pbackbuf ? -1 : (*--Pbackp=(c)))
18
19 void ungets( start, n )
20 char *start;
21 int n;
22 {
23     /* Push back the last n characters of the string by working backwards
24      * through the string.
25      */
26
27     char *p = start + strlen(start);      /* Find the end of the string. */
28
29     while( --p >= start && --n >= 0 )
30         if( unget(*p) == -1 )
31             fprintf( stderr, "Pushback-stack overflow\n" );
32 }
```

Input-system design criteria.

2.3.1 An Example Input System*

This section describes the input system used by a LEX-generated lexical analyzer (though not by LEX itself). This input system has many of the qualities that are desirable in a compiler's input system, and as such can be taken as characteristic. There are, of course, other solutions to the input problem,³ but the current one has proven quite workable and provides a good example of the sorts of problems that come up. Several design criteria must be met:

- The routines should be as fast as possible, with little or no copying of the input strings.
- Several characters of pushback and lookahead must be available.
- Lexemes of a reasonable length must be supported.
- Both the current and previous lexeme must be available.

3. A system that's more appropriate in a Pascal environment is described in [Aho] pp. 88–92.

- Disk access should be efficient.

To meet the last criterion, consider how a disk is accessed by most operating systems. All disks are organized into *sectors* which must be read as a unit. Disk reads must be performed in sector-sized chunks. If your disk has 512-byte sectors, then you must read 512 bytes at a time. This limitation is imposed by the hardware. If you request one byte from the disk, the operating system will read an entire sector into a buffer, and then return a single byte from the buffer. Subsequent reads get characters from the buffer until it is exhausted, and only then is a new sector read. Some operating systems (MS-DOS is a case in point) impose further constraints in that a group of several sectors (called a *cluster* or *block*) is the smallest possible unit that can be accessed directly, so you have to read an entire cluster at once. The minimum number of bytes that can be read is called an *allocation unit*, and, for the sake of efficiency, all reads must be done in terms of allocation units. That is, the number of bytes read from the disk at any one time should be a multiple of the allocation unit. Typically, the larger the buffer, the shorter the read time. Many operating systems reward you for doing block-sized transfers by not buffering the input themselves, as would be the case when an odd-sized block was requested from a low-level read call. The operating system transfers the data directly from the disk into your own buffer, thereby eliminating one level of copying and decreasing the read time. For example, MS-DOS read and write times improve dramatically when you read 32K bytes at a time.

The other design criteria are met by using a single input buffer and several pointers. My system is pictured in Figure 2.2. The drawing on the top shows the condition of the buffer just after it is loaded the first time. BUFSIZE is the actual buffer size. MAXLEX is the maximum lexeme length, and the disk reads are always in multiples of this number. Start_buf marks the physical start of the buffer, and END marks the physical end of the buffer. End_buf points at the logical end of buffer. (Since reads are in multiples of MAXLEX, and since the buffer itself isn't an even multiple of MAXLEX in length, there is usually a scrap of wasted space at the end of the buffer. End_buf points just past the last valid character in the buffer.) Finally, Next points at the next input character. (I'll discuss DANGER and MAXLOOK momentarily.)

The middle picture in Figure 2.2 shows the buffer in its normal state, after the lexical analyzer has processed several tokens. Various pointers have been set to mark the boundaries of various lexemes: pMark points at the beginning of the previous lexeme, sMark points at the beginning of the current lexeme, and eMark points at the end of the current lexeme. The lexical analyzer has scanned several characters past the end of the current lexeme (Next is to the right of eMark). If the lexical analyzer finds a longer lexeme than the current one, all it need do is move the eMark to the current input position. If, on the other hand, it finds that it has read too far in the input, it can push back all the extra characters by setting Next back to the eMark.

Returning to MAXLOOK, this constant is the number of lookahead characters that are supported. The DANGER marker tells the input routines when the Next pointer is getting too close to the end of the buffer (there must be at least MAXLOOK characters to the right of Next). When Next crosses the DANGER point, a buffer flush is triggered, giving us the situation shown in the bottom picture in Figure 2.2. All characters between the pMark and the last valid character (pointed to by End_buf) have been shifted to the far left of the buffer. The input routines fill the remainder of the buffer from the disk, reading in as many MAXLEX-sized chunks as will fit. The End_buf pointer is adjusted to mark the new end of the buffer, and DANGER scales automatically: it's positioned relative to the new end of buffer. This may seem like a lot of copying, but in practice the lexemes are not that large, especially in comparison to the buffer size. Consequently, flushes don't happen very often, and only a few characters are copied when they do happen.

Disk access, sectors.

Clusters, blocks.

Allocation units.

Input system organization: buffers and pointers.

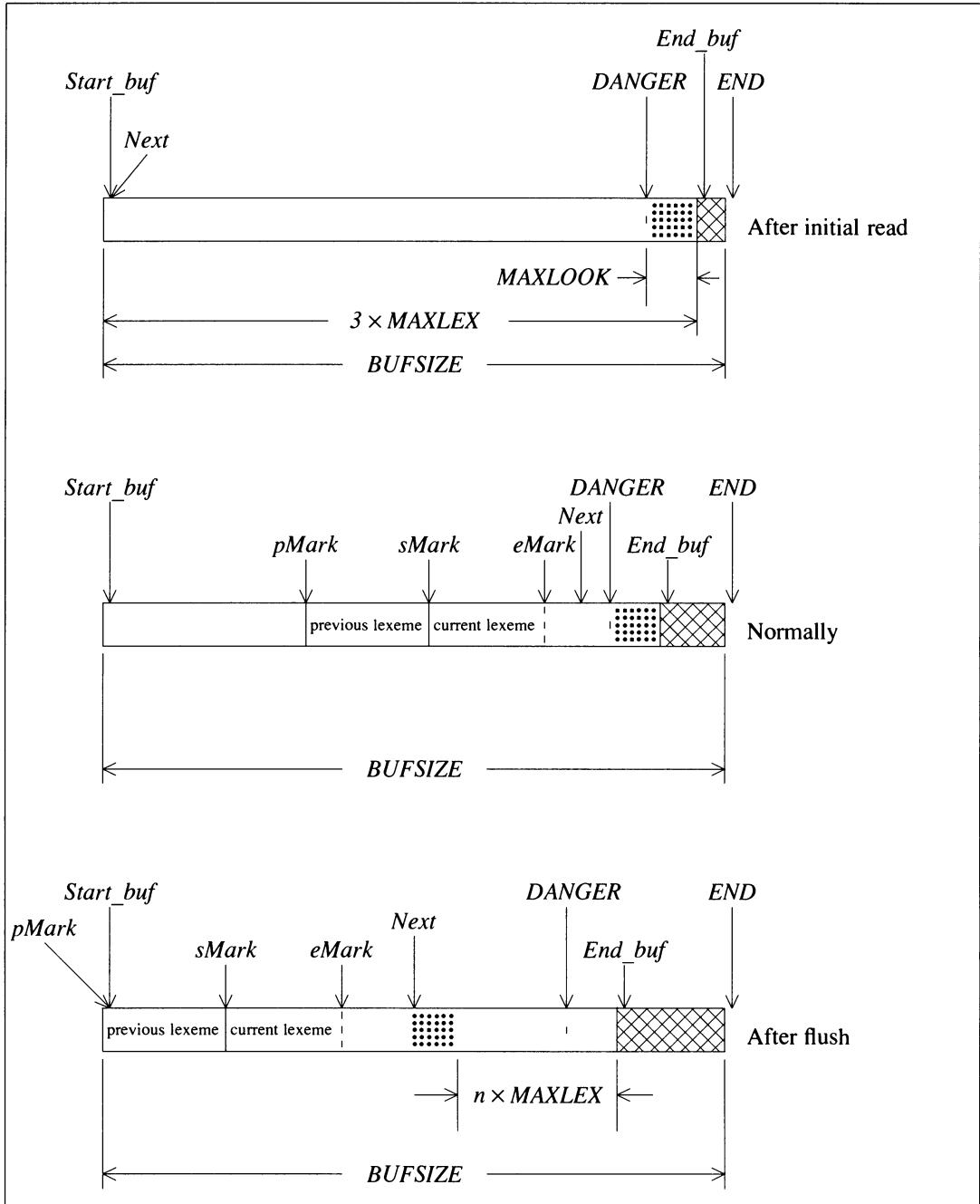
Buffer pointers, BUFSIZE, MAXLEX, END, Start_buf, End_buf.

Lexeme markers, pMark, sMark, eMark.

MAXLOOK, DANGER.

Buffer flush.

Figure 2.2. The Input Buffer



This approach has many advantages, the main one being a lack of copying. The lexical analyzer can just return the *sMark* as the pointer to the next lexeme, without having to copy it anywhere. Similarly a pushback is a single assignment (*Next*=*eMark*) rather than a series of pushes and pops. Finally, the disk reads themselves are reasonably efficient because they're done in block-sized chunks.

2.3.2 An Example Input System—Implementation

The foregoing is all implemented by the variables and macros declared at the top of *input.c*, in Listing 2.2. At this point we've looked at most of them, the others are discussed as they're used in the code. The macro definitions on lines 15 to 19 take care of a few system dependencies—`COPY()` is mapped to `memmove()` for the Microsoft C compiler [because Microsoft's `memcpy()` doesn't support overlapping strings].

Portability problems:
`COPY()`.

Listing 2.2. *input.c*—Macros and Data Structures

```

1  #include <stdio.h>
2  #include <stdlib.h>
3  #include <fcntl.h>
4  #include <tools/debug.h>
5
6  #include <tools/l.h>    /* Needed only for prototypes */
7  #include <string.h>      /* " */
8
9  /-----
10 * INPUT.C:   The input system used by LeX-generated lexical analyzers.
11 *-----
12 * System-dependent defines.
13 */
14
15 #ifdef MSDOS
16 #define COPY(d,s,a) memmove(d,s,a)
17 #else
18 #define COPY(d,s,a) memcpy(d,s,a)
19#endif
20
21 #define STDIN    0           /* standard input */
22
23 /-----
24
25 #define MAXLOOK     16          /* Maximum amount of lookahead */
26 #define MAXLEX      1024         /* Maximum lexeme sizes. */
27
28 #define BUFSIZE   ( (MAXLEX * 3) + (2 * MAXLOOK) )    /* Change the 3 only */
29
30 #define DANGER     ( End_buf - MAXLOOK )        /* Flush buffer when Next */
31                                /* passes this address */
32
33 #define END       (&Start_buf[BUFSIZE])        /* Just past last char in buf */
34
35 #define NO_MORE_CHARS() ( Eof_read && Next >= End_buf )
36
37 typedef unsigned char uchar;
38
39 PRIVATE uchar Start_buf[BUFSIZE]; /* Input buffer */
40 PRIVATE uchar *End_buf = END;    /* Just past last character */
41 PRIVATE uchar *Next = END;     /* Next input character */
42 PRIVATE uchar *sMark = END;    /* Start of current lexeme */
43 PRIVATE uchar *eMark = END;    /* End of current lexeme */
44 PRIVATE uchar *pMark = NULL;   /* Start of previous lexeme */
45 PRIVATE int pLineno = 0;        /* Line # of previous lexeme */
46 PRIVATE int pLength = 0;        /* Length of previous lexeme */
47
48 PRIVATE int Inp_file = STDIN;   /* Input file handle */
49 PRIVATE int Lineno = 1;         /* Current line number */
50 PRIVATE int Mline = 1;          /* Line # when mark_end() called */
51 PRIVATE int Termchar = 0;       /* Holds the character that was */

```



Listing 2.2. continued...

```

52                                     /* overwritten by a \0 when we      */
53                                     /* null terminated the last       */
54                                     /* lexeme.                      */
55 PRIVATE int    Eof_read   = 0;      /* End of file has been read.    */
56                                     /* It's possible for this to be  */
57                                     /* true and for characters to   */
58                                     /* still be in the input buffer. */
59
60 extern int open(), close(), read();
61
62 PRIVATE int  (*Openp) ()  = open ;  /* Pointer to open function      */
63 PRIVATE int  (*Closep) () = close ; /* Pointer to close function     */
64 PRIVATE int  (*Readp) ()  = read ;  /* Pointer to read function      */

```

Change low-level input routines, `ii_io()`.

Open new input file, `ii_newfile()`.

Reassigning standard input.

Binary (untranslated) input.

First read is delayed until first advance.

The actual code starts in Listing 2.3 with two initialization functions. The first, `ii_io()` on line 65 of Listing 2.3, is used to change the low-level input functions that are used to open files and fill the buffer. You may want to do this if you're getting input directly from the hardware, from a string, or doing something else that circumvents the normal input mechanism. This way you can use a L^EX-generated lexical analyzer in an unusual situation without having to rewrite the input system.

The `ii_newfile()` routine on line 81 of Listing 2.3 is the normal mechanism for opening a new input file. It is passed the file name and returns the file descriptor (not the FILE pointer) for the opened file, or -1 if the file couldn't be opened. The previous input file is closed unless it was standard input. `ii_newfile()` does not actually read the first buffer; rather, it sets up the various pointers so that the buffer is loaded the first time a character is requested. This way, programs that never call `ii_newfile()` will work successfully, getting input from standard input. The problem with this approach is that you must read at least one character before you can look ahead in the input (otherwise the buffer won't be initialized). If you need to look ahead before advancing, use:

```

ii_advance(); /* Read first bufferfull of input */
ii_pushback(1); /* but put back the first character */

```

The default input stream [used if `ii_newfile()` is never called] is standard input. You can reassign the input to standard input (say, after you get input from a file) by calling:

```
ii_newfile(NULL);
```

It's also okay to do a `ii_newfile("/dev/tty")` (in both MS-DOS and UNIX), but input is actually taken from the physical console in this case. Redirection won't work. An `ii_newfile(NULL)` allows for redirected input, however.

Note that the indirect `open()` call on line 103 of Listing 2.3 uses the `O_BINARY` input mode in MS-DOS systems (it's mapped to zero in UNIX systems). A CR-LF (carriage-return, linefeed) pair is not translated into a single '\n' when binary-mode input is active. This behavior is desirable in most L^EX applications, which treat both CR and LF as white space. There's no point wasting time doing the translation. The lack of translation might cause problems if you're looking for an explicit '\n' in the input, though.

Note that the input buffer is not read by `ii_newfile()`; rather, the various pointers are initialized to point at the end of the buffer on lines 111 to 114 of Listing 2.3. The actual input routine (`advance()`, discussed below) treats this situation the same as it would the Next pointer crossing the DANGER point. It shifts the buffer's tail all the way to the left (in this case the tail is empty so no characters are shifted, but the pointers are

Listing 2.3. *input.c*—Initialization Routines

```

65 void      ii_io( open_funct, close_funct, read_funct )
66 int       (*open_funct)();
67 int       (*close_funct)();
68 int       (*read_funct)();
69 {
70     /* This function lets you modify the open(), close(), and read() functions
71     * used by the i/o system. Your own routines must work like the real open,
72     * close, and read (at least in terms of the external interface. Open should
73     * return a number that can't be confused with standard input (not 0).
74     */
75
76     Openp = open_funct;
77     Closep = close_funct;
78     Readp = read_funct;
79 }
80 /*-----*/
81 int      ii_newfile( name )
82 char    *name;
83 {
84     /* Prepare a new input file for reading. If newfile() isn't called before
85     * input() or input_line() then stdin is used. The current input file is
86     * closed after successfully opening the new one (but stdin isn't closed).
87     */
88
89     * Return -1 if the file can't be opened; otherwise, return the file
90     * descriptor returned from open(). Note that the old input file won't be
91     * closed unless the new file is opened successfully. The error code (errno)
92     * generated by the bad open() will still be valid, so you can call perror()
93     * to find out what went wrong if you like. At least one free file
94     * descriptor must be available when newfile() is called. Note in the open
95     * call that O_BINARY, which is needed in MS-DOS applications, is mapped
96     * to 0 under UNIX (with a define in <tools/debug.h>).
97     */
98
99     int fd;                      /* File descriptor */
100
101    MS( if( strcmp(name, "/dev/tty") == 0 ) )
102    MS(         name = "CON" );
103
104    if( (fd = !name ? STDIN : (*Openp)(name, O_RDONLY|O_BINARY)) != -1 )
105    {
106        if( Inp_file != STDIN )
107            (*Closep)( Inp_file );
108
109        Inp_file = fd;
110        Eof_read = 0;
111
112        Next      = END;
113        sMark     = END;
114        eMark     = END;
115        End_buf   = END;
116        Linenos  = 1;
117        Mline     = 1;
118    }
119    return fd;
120 }
```

moved), and then loads the buffer from the disk. I've taken this approach because it's sometimes convenient to open a default input file at the top of a program, which is then overridden by a command-line switch or the equivalent later on in the same program. There's no point in reading from a file that's not going to be used, so the initial read is delayed until a character is requested.

Access functions.

Limiting scope,
PRIVATE.

Functions to access lexemes: `ii_text()`,
`ii_length()`,
`ii_lineno()`,
`ii_ptext()`,
`ii_plength()`,
`ii_plinen()`.

Functions to mark lexeme boundaries,
`ii_mark_start()`,
`ii_mark_end()`.

Move start marker,
`ii_move_start()`.

Restore pointer to previous mark,
`ii_to_mark()`.
 Mark previous lexeme,
`ii_mark_prev()`.

Advance input pointer,
`ii_advance()`.

The `input.c` file continues in Listing 2.4 with several small *access* functions. For maintenance reasons, it is desirable to limit external access of global variables, because the linker assumes that two global variables with the same name are the same variable. If you inadvertently declare two variables with the same name, one of them will seem to magically change its value when a subroutine that accesses the other is called. You can avoid this problem by declaring the variables `static`, thereby limiting their scope to the current file. `PRIVATE` is mapped to `static` in `debug.h`, discussed in Appendix A.

It's still necessary for external subroutines to access these variables however, and the safest way to do so is through the small routines in Listing 2.4. These subroutines are used for maintenance reasons only—two subroutines with the same name will result in an error message from the linker, unlike two variables with the same name, which are silently merged.

The `ii_text()`, `ii_length()`, and `ii_lineno()` routines (lines 120 to 122 of Listing 2.4) return a pointer to the current lexeme, the lexeme's length, and the line number for the last character in the lexeme. The `ii_ptext()`, `ii_plength()`, and `ii_plinen()` routines (lines 123 to 125) do the same thing, but for the previous lexeme. The `ii_mark_start()` routine (line 127) moves the `sMark` to the current input position (pointed to by `Next`). It also makes sure that the end-of-lexeme marker (`eMark`) is not to the left of the start marker. `ii_mark_end()` (line 134) does the same for the end marker (`eMark`). It also saves the current line number in `Mline`, because the lexical analyzer might sweep past a newline when it scans forward looking for a new lexeme. The input line number must be restored to the condition it was in before the extra newline was scanned when the analyzer returns to the previous end marker.

The `ii_move_start()` routine on line 140 of Listing 2.4 lets you move the start marker one space to the right. It returns the new start marker on success, `NULL` if you tried to move past the end marker (`sMark` is not modified in this last case). `ii_to_mark()` (line 148) restores the input pointer to the last end mark. Finally, `ii_mark_prev()` modifies the previous-lexeme marker to reference the same lexeme as the current-lexeme marker. Typically, `ii_mark_prev()` is called by the lexical analyzer just before calling `ii_mark_start()` (that is, just before it begins to search for the next lexeme).

The next group of subroutines, in Listings 2.5 and 2.6, comprise the advance and buffer-flush functions. `ii_advance()`, on line 168 of Listing 2.5, returns the next input character and advances past it. The code on lines 180 to 191 is provided for those situations where you want an extra newline appended to the beginning of a file. L^EX needs this capability for processing the start-of-line anchor—a mechanism for recognizing strings only if they appear at the far left of a line. Such strings must be preceded by a newline, so an extra newline has to be appended in front of the first line of the file; otherwise, the anchored expression wouldn't be recognized on the first line.⁴

4. L^EX Usage Note: This pushback could conceivably cause problems if there is no regular expression in the L^EX input file to absorb the newline, and YYBADINP is also #defined (You'll get an error message in this case). A regular expression that absorbs white space is usually present, however.

Listing 2.4. *input.c*—Small Access Routines and Marker Movement

```

120 PUBLIC char *ii_text() { return( sMark ); }
121 PUBLIC int ii_length() { return( eMark - sMark ); }
122 PUBLIC int ii_lineno() { return( Lineno ); }
123 PUBLIC char *ii_ptext() { return( pMark ); }
124 PUBLIC int ii_plength() { return( pLength ); }
125 PUBLIC int ii_plinenno() { return( pLineno ); }

126
127 char *ii_mark_start()
128 {
129     Mline = Lineno;
130     eMark = sMark = Next;
131     return( sMark );
132 }
133
134 PUBLIC char *ii_mark_end()
135 {
136     Mline = Lineno ;
137     return( eMark = Next );
138 }
139
140 PUBLIC char *ii_move_start()
141 {
142     if ( sMark >= eMark )
143         return NULL;
144     else
145         return ++sMark ;
146 }
147
148 PUBLIC char *ii_to_mark()
149 {
150     Lineno = Mline ;
151     return( Next = eMark );
152 }
153
154 char *ii_mark_prev()
155 {
156     /* Set the pMark. Be careful with this routine. A buffer flush won't go past
157      * pMark so, once you've set it, you must move it every time you move sMark.
158      * I'm not doing this automatically because I might want to remember the
159      * token before last rather than the last one. If ii_mark_prev() is never
160      * called, pMark is just ignored and you don't have to worry about it.
161      */
162
163     pMark = sMark;
164     pLineno = Lineno;
165     pLength = eMark - sMark;
166     return( pMark );
167 }
```

The `NO_MORE_CHARS()` macro is used on line 193 to detect end of file. It was defined in the header as follows

Detect end of file,
`NO_MORE_CHARS()`.

```
#define NO_MORE_CHARS() ( Eof_read && Next >= End_buf )
```

`Eof_read` is set to true when end of file is encountered. You must use both `Eof_read` and `Next` to detect end of input because EOF might have been read while the lexical analyzer was looking ahead. In this case, characters may have been pushed back after

`Eof_read`.

End of input.

reading the EOF. You have to see both if end of file has been encountered and if the buffer is empty. This is a case where *end of input* and *end of file* are different things, because there still may be characters in the input buffer long after end of file has been read. The `ii_flush()` call on line 196 flushes the buffer if necessary, and the line number is advanced on line 199. The next input character is returned normally, 0 is returned on end of file, and -1 is returned if the buffer couldn't be flushed for some reason.

Listing 2.5. *input.c*—The Advance Function

```

168 int ii_advance()
169 {
170     /* ii_advance() is the real input function. It returns the next character
171      * from input and advances past it. The buffer is flushed if the current
172      * character is within MAXLOOK characters of the end of the buffer. 0 is
173      * returned at end of file. -1 is returned if the buffer can't be flushed
174      * because it's too full. In this case you can call ii_flush(1) to do a
175      * buffer flush but you'll lose the current lexeme as a consequence.
176      */
177
178     static int been_called = 0;
179
180     if( !been_called )
181     {
182         /* Push a newline into the empty buffer so that the LeX start-of-line
183          * anchor will work on the first input line.
184          */
185
186         Next = sMark = eMark = END - 1;
187         *Next = '\n';
188         --Lineno ;
189         --Mline ;
190         been_called = 1;
191     }
192
193     if( NO_MORE_CHARS() )
194         return 0;
195
196     if( !Eof_read && ii_flush(0) < 0 )
197         return -1;
198
199     if( *Next == '\n' )
200         Lineno++;
201
202     return( *Next++ );
203 }
```

Flush input buffer,
`ii_flush()`.

force.

The actual buffer flush is done by `ii_flush()`, which starts at the top of Listing 2.6. The test on line 248 checks to see that there will be enough room after the move to load a new MAXLEX-sized bufferfull of characters—there might not be if the buffer contains two abnormally long lexemes. The test evaluates true if there isn't enough room. Normally, the routine returns -1 if there's no room, and 1 is returned if everything is okay. If the `force` argument is true, however, the buffer is flushed even if there's no room, and 1 is returned. The flush is forced by setting the start marker to the current input position and the `left_edge` of the character to be shifted to the `Next` pointer, effectively destroying the current lexeme. The code on lines 259 and 246 figures out

how many characters have to be copied (`copy_amt`) and the distance that they have to be moved (`shift_amt`). The shift is done on line 260, and a new buffer is loaded by the `ii_fillbuf()` call on line 262. `COPY` was defined earlier (on line 16 of Listing 2.2) to map to either `memmove()` or `memcpy()`, depending on the compilation environment. The rest of the routine adjusts the various markers to compensate for the move.

`copy_amt, shift_amt.`

Listing 2.6. *input.c*—Buffer Flushing

```

204 int      ii_flush( force )
205 int      force;
206 {
207     /* Flush the input buffer. Do nothing if the current input character isn't
208     * in the danger zone, otherwise move all unread characters to the left end
209     * of the buffer and fill the remainder of the buffer. Note that input()
210     * flushes the buffer willy-nilly if you read past the end of buffer.
211     * Similarly, input_line() flushes the buffer at the beginning of each line.
212     *
213     *          pMark      DANGER
214     *          |          |
215     *          Start_buf      sMark  eMark |Next   End_buf
216     *          |          ||      |  ||    |
217     *          V          VV    V   VV   V
218     *          +-----+-----+-----+
219     *          | this is already read | to be done yet | waste |
220     *          +-----+-----+-----+
221     *          |          |          |          |
222     *          |<---- shift_amt ---->|<-- copy_amt -->|          |
223     *          |
224     *          |<----- BUFSIZE ----->|
225     *
226     * Either the pMark or sMark (whichever is smaller) is used as the leftmost
227     * edge of the buffer. None of the text to the right of the mark will be
228     * lost. Return 1 if everything's ok, -1 if the buffer is so full that it
229     * can't be flushed. 0 if we're at end of file. If "force" is true, a buffer
230     * flush is forced and the characters already in it are discarded. Don't
231     * call this function on a buffer that's been terminated by ii_term().
232 */
233
234     int      copy_amt, shift_amt ;
235     uchar *left_edge;
236
237     if( NO_MORE_CHARS() )
238         return 0;
239
240     if( Eof_read )           /* nothing more to read */
241         return 1;
242
243     if( Next >= DANGER || force )
244     {
245         left_edge = pMark ? min(sMark, pMark) : sMark;
246         shift_amt = left_edge - Start_buf ;
247
248         if( shift_amt < MAXLEX )           /* if(not enough room) */
249         {
250             if( !force )
251                 return -1;
252
253             left_edge = ii_mark_start(); /* Reset start to current character */
254             ii_mark_prev();
255

```



Listing 2.6. continued...

```

256             shift_amt = left_edge - Start_buf ;
257         }
258
259         copy_amt = End_buf - left_edge;
260         COPY( Start_buf, left_edge, copy_amt );
261
262         if( !ii_fillbuf( Start_buf + copy_amt ) )
263             ferr("INTERNAL ERROR, ii_flush: Buffer full, can't read.\n");
264
265         if( pMark )
266             pMark -= shift_amt;
267
268         sMark -= shift_amt;
269         eMark -= shift_amt;
270         Next -= shift_amt;
271     }
272
273     return 1;
274 }
275
276 /*-----*/
277
278 PRIVATE int ii_fillbuf( starting_at )
279 unsigned char *starting_at;
280 {
281     /* Fill the input buffer from starting_at to the end of the buffer.
282      * The input file is not closed when EOF is reached. Buffers are read
283      * in units of MAXLEX characters; it's an error if that many characters
284      * cannot be read (0 is returned in this case). For example, if MAXLEX
285      * is 1024, then 1024 characters will be read at a time. The number of
286      * characters read is returned. Eof_read is true as soon as the last
287      * buffer is read.
288      *
289      * PORTABILITY NOTE: I'm assuming that the read function actually returns
290      *                   the number of characters loaded into the buffer, and
291      *                   that that number will be < need only when the last chunk of the file is
292      *                   read. It's possible for read() to always return fewer than the number of
293      *                   requested characters in MS-DOS untranslated-input mode, however (if the
294      *                   file is opened without the O_BINARY flag). That's not a problem here
295      *                   because the file is opened in binary mode, but it could cause problems
296      *                   if you change from binary to text mode at some point.
297      */
298
299     register unsigned need,          /* Number of bytes required from input. */
300             got;           /* Number of bytes actually read. */
301
302     need = ((END - starting_at) / MAXLEX) * MAXLEX ;
303
304     D( printf( "Reading %d bytes\n", need );    )
305
306     if( need < 0 )
307         ferr("INTERNAL ERROR (ii_fillbuf): Bad read-request starting addr.\n");
308
309     if( need == 0 )
310         return 0;
311
312     if( (got = (*Readp)(Inp_file, starting_at, need)) == -1 )
313         ferr("Can't read input file\n");
314

```



Listing 2.6. continued..

```

315     End_buf = starting_at + got ;
316
317     if( got < need )
318         Eof_read = 1;           /* At end of file */
319
320     return got;
321 }
```

The final routine in Listing 2.6 is `ii_fillbuf()`, starting on line 278. It is passed a base address, and loads as many MAXLEX-sized chunks into the buffer as will fit. The `need` variable, initialized on line 302, is the amount needed. The logical-end-of-buffer marker is adjusted on line 315. Note that a single `read()` call does the actual read on line 312. (`Readp` is initialized to point at `read()` when it is declared up at the top of the file.) This can cause problems when a lexeme can span a line, and input is fetched from a line-buffered input device (such as the console). You'll have to use `ii_io()` to supply an alternate read function, in this case.

Listing 2.7 shows the lookahead function, `ii_look()`. It returns the character at the offset from the current character that's specified in its argument. An `ii_look(0)` returns the character that was returned by the most recent `ii_advance()` call, `ii_look(1)` is the following character, `ii_look(-1)` is the character that precedes the current one. MAXLOOK characters of lookahead are guaranteed, though fewer might be available if you're close to end of file. Similarly, lookback (with a negative offset) is only guaranteed as far as the start of the buffer (the `pMark` or `sMark`, whichever is smaller). Zero is returned if you try to look past end or start of the buffer, EOF if you try to look past end of file.

Load input buffer,
`ii_fillbuf()`.
need.

Lookahead, `ii_look()`.

Listing 2.7. *input.c*—Lookahead

```

322 int      ii_look( n )
323 {
324     /* Return the nth character of lookahead, EOF if you try to look past
325      * end of file, or 0 if you try to look past either end of the buffer.
326      */
327
328     uchar      *p;
329
330     p = Next + (n-1) ;
331
332     if( Eof_read && p >= End_buf )
333         return EOF;
334
335     return( p < Start_buf || p >= End_buf ) ? 0 : *p ;
336 }
```

Listing 2.8 contains the pushback function, `ii_pushback(n)`. It is passed the number of characters to push back. For example, `ii_pushback(5)` pushes back the five most recently read characters. If you try to push past the `sMark`, only the characters as far as the `sMark` are pushed and 0 is returned (1 is returned on a successful push). If you push past the `eMark`, the `eMark` is moved back to match the current character. Unlike `ungetc()`, you can indeed push back characters after EOF has been reached.

`ii_pushback(n)`.

The remainder of `input.c`, in Listing 2.9, provides support for '\0'-terminated strings. These routines are not—strictly speaking—necessary, because the lexeme

\0-terminated-string support.

Listing 2.8. *input.c*—Pushback

```

337 int ii_pushback( n )
338 {
339     /* Push n characters back into the input. You can't push past the current
340      * sMark. You can, however, push back characters after end of file has
341      * been encountered.
342      */
343
344     while( --n >= 0 && Next > sMark )
345     {
346         if( *--Next == '\n' || !*Next )
347             --Lineno;
348     }
349
350     if( Next < eMark )
351     {
352         eMark = Next;
353         Mline = Lineno;
354     }
355
356     return( Next > sMark );
357 }
```

length is always available. It's occasionally useful to have a terminator on the string, however. Note that these functions should be used exclusively after the string has been terminated—the other input functions will not work properly in this case.

Terminate/unterminate
current lexeme,
ii_term(),
ii_unterm().

The termination is done with a call to *ii_term()* (on line 358). It saves the character pointed to by *Next* in a variable called *Termchar*, and then overwrites the character with a '\0'. The *ii_unterm()* function (on line 366) puts everything back.

Listing 2.9. *input.c*—Support for '\0'-terminated Strings

```

358 void ii_term()
359 {
360     Termchar = *Next ;
361     *Next     = '\0' ;
362 }
363
364 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
365
366 void ii_unterm()
367 {
368     if( Termchar )
369     {
370         *Next = Termchar;
371         Termchar = 0;
372     }
373 }
374
375 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
376
377 int ii_input()
378 {
379     int rval;
380 }
```



Listing 2.9. continued...

```

382     if( Termchar )
383     {
384         ii_unterm();
385         rval = ii_advance();
386         ii_mark_end();
387         ii_term();
388     }
389     else
390     {
391         rval = ii_advance();
392         ii_mark_end();
393     }
394
395     return rval;
396 }
397
398 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
399
400 void    ii_unput( c )
401 {
402     if( Termchar )
403     {
404         ii_unterm();
405         if( ii_pushback(1) )
406             *Next = c;
407         ii_term();
408     }
409     else
410     {
411         if( ii_pushback(1) )
412             *Next = c;
413     }
414 }
415
416 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
417
418 int     ii_lookahead( n )
419 {
420     return (n == 1 && Termchar) ? Termchar : ii_look(n) ;
421 }
422
423 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
424
425 int     ii_flushbuf()
426 {
427     if( Termchar )
428         ii_unterm();
429
430     return ii_flush(1);
431 }
```

This approach is better than putting the `ii_unterm()` code into `ii_advance()`, because the latter approach slows down all `ii_advance()` calls. On the other hand, you have to remember to call `ii_unterm()` before calling `ii_advance()`. For this reason, an `ii_input()` function has been provided (on line 377) to make sure that the lexeme is unterminated and then retermined correctly. That is, `ii_input()` is a well-behaved input function meant to be used directly by the user. The function also

`ii_unput()`.

`ii_lookahead()`.

Ring buffers.

Why ring buffers are inappropriate here.

Lookup tables in hard-coded scanners.

Hard-coded scanners:
advantages and disadvantages.

moves the end marker, making the lexeme one character longer (moving the null terminator if necessary), and it returns the new input character, or 0 at end of file. -1 is returned if another character couldn't be read because the buffer was full.

`ii_unput()` (on line 400) is a reverse-input function. It backs up the input one notch and then overwrites the character at that position with its argument. `ii_unput()` works correctly on both terminated and unterminated buffers, unlike `ii_pushback()`, which can't handle the terminator.

The `ii_lookahead()` function bears the same relation to `ii_look()` that `ii_input()` bears to `ii_advance()`. That is, `ii_lookahead(1)` functions correctly for strings that have been terminated with `ii_term()` calls, `ii_look()` does not. Similarly, `ii_flushbuf()` flushes a terminated buffer by unterminating it before calling `ii_flush()`.

One final note on strategy. The buffer-flush approach that I've used here allows me to take advantage of C's pointer mechanism when scanning the input. This approach isn't appropriate in a language like FORTRAN, where arrays must be referenced using an index. Here, you're better off using a *circular array* or *ring buffer*. For example, the input buffer would be declared with

```
char input_buf[ SIZE ];
```

and the next character would be accessed with

```
x = input_buf[ current_character % SIZE ];
```

You would load a new chunk from the far left of the array when `current_character` was greater than or equal to `SIZE`, being careful not to overwrite the current lexeme in the process.

The problem here is that a lexeme can span the buffer. That is, a situation may arise where the first half of a lexeme is far right of `input_buf` and the other half is at the far left. As long as you're accessing all the characters with an array index modulus the array size, this is not a problem. C, however, wants its strings in contiguous memory so that it can scan through them using a pointer. Moreover, the array index and modulus operation needed to access every character is inherently less efficient than a simple pointer access; more inefficient, even, than the moves that are part of a buffer flush. Consequently, a ring buffer isn't particularly appropriate in a C implementation.

2.4 Lexical Analysis*

Now that we've developed a set of input routines, we need to apply them in a lexical-analysis application. There are two approaches to lexical analysis, both useful. First, you can hard code the analyzer, recognizing lexemes with nested `if/else` statements, switches, and so forth. If the lexemes aren't too long one effective approach uses a series of lookup tables to recognize tokens. (Lookup tables tend to be faster than `switchs` or `if/else` statements.) Listing 2.10 shows such a system for recognizing the following tokens:

```
>  >=  <  <=  ==  !=
```

The basic strategy is to vector from one table to another until a complete lexeme is identified.

The hard-coded approach has its advantages—hard-coded lexical analyzers tend to be very efficient, but hard-coded analyzers are difficult to maintain. When you're developing a new language, it's handy to be able to add new tokens to the language or take some away without too much work. This problem is solved by programs like `lex`

Listing 2.10. Using Lookup Tables for Character Recognition

```

1 #define LESS_THAN 1
2 #define GREATER_THAN 2
3 #define EQUAL 3
4 #define NOT 4
5 #define LESS_THAN_OR_EQUAL 5
6 #define GREATER_THAN_OR_EQUAL 6
7 #define NOT_OR_EQUAL 7
8 #define ASSIGN 8
9
10 #define ERROR -1
11 #define CONTINUE 0
12
13 #define SIZE_OF_CHARACTER_SET 128
14
15 char first [ SIZE_OF_CHARACTER_SET ];
16 char second[ SIZE_OF_CHARACTER_SET ];
17 int s, f;
18 ...
19 memset( first, -1, SIZE_OF_CHARACTER_SET ); /* Initialize to error token. */
20 memset( second, -1, SIZE_OF_CHARACTER_SET ); /* Note that there's an implicit */
21 /* conversion of -1 to 255 here */
22 /* (signed int to unsigned char).*/
23 ...
24 first [ '>' ] = GREATER_THAN;
25 first [ '<' ] = LESS_THAN;
26 first [ '!' ] = NOT;
27 first [ '=' ] = ASSIGN;
28 second[ '=' ] = EQUAL;
29
30 c = getchar();
31
32 if( (f = first[c]) == ERROR ) /* discard bad character */
33     return ERROR;
34
35 if( (s = second[c]) == ERROR ) /* 1-character lexeme */
36 {
37     ungetchar();
38     return( f );
39 }
40 else /* 2-character lexeme */
41 {
42     if( s == EQUAL )
43         switch( f )
44         {
45             case ASSIGN:      return EQUAL ;
46             case LESS_THAN:  return LESS_THAN_OR_EQUAL ;
47             case GREATER_THAN: return GREATER_THAN_OR_EQUAL ;
48             case NOT:        return NOT_EQUAL ;
49         }
50     return ERROR; /* discard both characters */
51 }
```

and L^EX, which translate a description of a token set into a table-driven lexical analyzer. (Hereafter, when I say L^EX, I'm actually referring to both programs). L^EX itself just creates the tables, the remainder of the analyzer is the same for all L^EX-generated source code. The fact that a L^EX-generated analyzer is typically slower than a hard-coded one is, more often than not, a small price to pay for faster development time. Once the

language is stable, you can always go back and hard code a lexical analyzer if it's really necessary to speed up the front end.

2.4.1 Languages*

Alphabets, strings, words.

ϵ , the empty string.

Empty versus null strings

Languages, sentences, grammars.

Before looking at L^EX itself, we'll need a little theory. First some definitions. An *alphabet* is any finite set of symbols. For example, the ASCII character set is an alphabet; the set $\{'0', '1'\}$ is a more-restricted, binary alphabet. A *string* or *word* is a sequence of alphabetic symbols. In practical terms, a string is an array of characters. There is also the special case of an empty string, represented by the symbol ϵ (pronounced epsilon). In C, the '`\0`' is not part of the input alphabet. As a consequence, it can be used as an end-of-string marker because it cannot be confused with any of the characters in the string itself, all of which are part of the input alphabet. An empty string in C can then be represented by an array containing a single '`\0`' character. Note, here, that there's an important difference between ϵ , an empty string, and a *null* string. The former is an array containing the end-of-string marker. The latter is represented by a NULL pointer—a pointer that doesn't point anywhere. In other words, there is no array associated with a null string.

A *language* is a set of strings that can be formed from the input alphabet. A *sentence* is a sequence of the strings that comprise a language. A language can be as small as one string and still be useful. (Zero-element languages are possible, but not of much utility.) The ordering of strings within the sentence is defined by a collection of syntactic rules called a *grammar*. Note that this definition does not attribute meaning to any of the strings and this limitation has important practical consequences. The lexical analyzer doesn't understand meaning. It has to distinguish tokens solely on the basis of surrounding context—by looking at the characters that surround the current word, without regard to the syntactic or semantic structure of the input sentence (the tokens that precede and follow the current token). [Aho] introduces several other useful terms and definitions, paraphrased here:⁵

prefix	A <i>prefix</i> is a string composed of the characters remaining after zero or more symbols have been deleted from the end of a string: "in" is a prefix of "inconsequential". Officially, ϵ is a prefix of every string.
suffix	A <i>suffix</i> is a string formed by deleting zero or more symbols from the front of a string. "ible" is a suffix of "incomprehensible". The suffix is what's left after you've removed a prefix. ϵ is a suffix of every string.
substring	A <i>substring</i> is what's left when you remove both a suffix and prefix: "age" is a substring of "unmanageable". Note that suffixes and prefixes are substrings (but not the other way around). Also ϵ , the empty string, is a substring of every string.
proper X	A <i>proper prefix</i> , <i>suffix</i> , or <i>substring</i> of the string x has at least one element and it is not the same as x . That is, it can't be ϵ , and it can't be identical to the original string.
sub-sequence	A <i>sub-sequence</i> of a string is formed by deleting zero or more symbols from the string. The symbols don't have to be contiguous, so "iiii" and "ssss" are both sub-sequences of "Mississippi".

5. [Aho], p. 93.

Several useful operations can be performed on strings. The *concatenation* of two strings is formed by appending all characters of one string to the end of another string. The concatenation of "fire" and "water" is "firewater". The empty string, ϵ , can be concatenated to any other string without modifying it. (In set theory, ϵ is the *identity element* for the concatenation operation. An arithmetic analogy is found in multiplication: 1 is the identity element for multiplication because $x \equiv x \times 1$.) The concatenation operation is sometimes specified with an operator (typically \cdot or \circ), so you can say that

$$\text{fire} \cdot \text{water} \equiv \text{firewater}$$

String concatenation.

 ϵ is the identity element.

If you look at concatenation as a sort of multiplication, then exponentiation makes sense. An expression like x^n represents x , repeated n times. You could define a language consisting of the eight legal octal digits with the following:

$$L(\text{octal}) = \{ 0, 1, 2, 3, 4, 5, 6, 7 \}$$

String exponentiation.

and then you could specify a three-digit octal number with $L(\text{octal})^3$.

Kleene Closure (*).

The exponentiation process can be generalized into the *closure* operations. If L is a language, then the *Kleene closure* of L is L repeated zero or more times. This operation is usually represented as L^* . In the case of a language comprised of a single character, L^* is that character repeated zero or more times. If the language elements are strings rather than single characters, L^* are the strings repeated zero or more times. For example, $L(\text{octal})^*$ is zero or more octal digits. If $L(v1)$ is a language comprised of the string Va and $L(v2)$ is a language comprised of the string $Voom$, then

$$L(v1)^* \cdot L(v2)$$

describes all of the following strings:

Voom VaVoom VaVaVoom VaVaVaVoom etc.

Positive closure (+).

The *positive closure* of L is L repeated one or more times, usually denoted L^+ . It's otherwise just like Kleene closure.

Set operations on languages, union (\cup).

Since languages are sets of symbols, most of the standard set operations can be applied to them. The most useful of these is union, denoted with the \cup operator. For example, if *letters* is a language containing all 26 letters [denoted by $L(\text{letters})$] and *digits* is a set containing all 10 digits [denoted by $L(\text{digits})$], then $\{L(\text{letters}) \cup L(\text{digits})\}$ is the set of alphanumeric characters. Union is the equivalent of a logical OR operator. (If x is an element of $A \cup B$, then it is a member of either A OR B .) Other set operations (like intersection) are, of course possible, but have less practical application.

The foregoing can all be applied to build a language from an alphabet and to define large languages (such as token sets) in terms of smaller languages (letters, digits, and so forth). For example

$$L(\text{digit}) = \{ 1, 2, 3, 4, 5, 6, 7, 8, 9 \}$$

$$L(\text{alpha}) = \{ a, b, c, \dots, z \}$$

you can say:

$$L(\text{digit})^+$$

is a decimal constant in C (one or more digits).

$$L(\text{digit})^*$$

is an optional decimal constant (zero or more digits).

$$L(\text{alpha}) \cup L(\text{digit})$$

is the set of alphanumeric characters.

$$(L(\text{alpha}) \cup L(\text{digit}))^*$$

is any number of alphanumeric characters.

$$L(\text{alpha}) \cdot (L(\text{alpha}) \cup L(\text{digit}))^*$$

is a C identifier.

2.4.2 Regular Expressions*

Programs like L^EX use the foregoing language theory to specify a token set for a lexical analyzer. The possible lexemes that correspond to individual tokens are all defined using a series of set operations applied to previously defined languages, with a base alphabet of the ASCII character set. The programs then translate that language specification into the C source code for a computer program that recognizes strings in the language.

Forming regular expressions. Metacharacters.

Both programs use a notation called *regular expressions* for this purpose. Strictly speaking, a regular expression is any well-formed formula over union, concatenation and Kleene closure—as was the case with the examples in the previous section. A practical implementation of regular expressions usually add other operations, however, to make them easier to use. I'll examine an extended regular-expression syntax in the current section.

The simplest regular expression is just a series of letters that match a sequence of the same letters in the input. Several special characters, called *metacharacters*, can be used to describe more complex strings. Though there are variations in the notation used for regular expressions, the following rules are used by L^EX to form a regular expression and can be taken as characteristic.⁶

Regular expression concatenation.

Wildcard character.

Start-of-line anchor.

End-of-line anchor.

Character classes.

- c *A single character that is not a metacharacter is a regular expression.* The character c forms a regular expression that matches the single character c.
 - ee *Two regular expressions concatenated form a regular expression that recognizes a match of the first expression followed by a match of the second.* If a, n, and d are regular expressions recognizing the characters a, n, and d, they can be concatenated to form the expression and which matches the pattern and in the input. Note that there's no explicit concatenation operator here, the two strings are just placed next to each other.
 - .
 - ^ *An up arrow anchors the pattern to the start of the line.* The pattern ^ and matches the string and only if it comprises the first three characters on the line (no preceding white space). Note that any newline character that precedes the and is not matched. That is, the newline is not part of the lexeme, even though its presence (or a start-of-file marker) is required for a successful match.
 - \$ *A dollar sign anchors the pattern to end of line.* The pattern and \$ matches the string only if it is the last three characters on the line (no following white space). Again, the newline character is not part of the lexeme. The pattern ^ and \$ matches the word only if it's the only thing on a line.
 - [...] [^...]
- Brackets match any of the characters enclosed in the brackets.* The [and] metacharacter form a *character class* which matches any of the characters listed. For example, [0123456789] matches any single decimal digit. Ranges of characters can be abbreviated using a dash, so [0-9] also matches a single decimal digit. [0-9A-Fa-f] matches a hexadecimal digit. [a-zA-Z] matches an alphabetic character. If the first character following the bracket is an up arrow (^), a *negative character class* (which matches any character except the ones specified) is formed. [^a-z]

6. Other UNIX utilities, like grep, vi, and sed, use a subset of these rules.

matches any character except a lower-case, alphabetic character. Only seven characters have special meaning inside a character class:

- { Start of macro name.
- } End of macro name.
-] End of character class.
- Range of characters.
- ^ Indicates negative character class.
- " Takes away special meaning of characters up to next quote mark.
- \ Takes away special meaning of next character.

Use \], \-, \\, and so forth, to put these into a class. Since other metacharacters such as *, ?, and + are not special here, the expression [*?+] matches a star, question mark, or plus sign. Also, a negative character class does not match a newline character. That is, [^a-z] actually matches anything except a lower-case character or newline. Note that a negative character class must match a character. That is, [^a-z]\$ does not match an empty line. The line must have at least one character, though it may not end in a nonalphanumeric character.

* + ?

*A regular expression followed by a * (pronounced star) matches that expression repeated zero or more times; a + matches one or more repetitions, a ? matches zero or one repetitions.* These three metacharacters represent *closure* operations. They are higher precedence than concatenation. ll?ama matches two strings: llama and lama. The expression l+ama matches lama, llama, and llllllllllllama. The expression l*ama matches all of the above, but also matches ama. The expression -0[xX][0-9a-fA-f]+ matches a hexadecimal number in C syntax; [0-7][0-7]? matches one or two octal digits.

Closure operators.

e{n,m}

Matches n to m repetitions of the expression e. This operator is recognized by **lex**, but not **LEX**.

Multiple matches.

e|e

Two regular expressions separated by a vertical bar recognize a match of the first expression OR a match of the second. OR is lower precedence than concatenation. The expression either|or matches either either or or.

The OR operator.

(e)

Parentheses are used for grouping. The expression:

Grouping .

```
(frank|john) ie
```

matches both frankie, and johnie. The expression

```
(frank| john) (ie)?
```

matches frank and john as well. You can add a newline to the characters recognized by a negative character class with something like this:

```
([^a-z]|\\en)
```

Add '0 to negative character class.

Surrounding a string that contains metacharacters with double quotes ("") or preceding a single metacharacter with a backslash (*) takes away its special meaning. (A character preceded by a backslash is said to be *escaped*.) The operator precedence is summarized in the Table 2.1. All operators associate left to right.

Escaping metacharacters, quotes marks.

Note that regular expressions can only define sequences of characters. They cannot do things like recognize any number of properly nested parentheses, something that can be recognized grammatically (by the parser). This is one of the main reasons that the lexical analyzer and parser are separate modules. The lexical analyzer is in charge of recognizing simple sequences of characters, and the parser recognizes more complex combinations.

Limitations of regular expressions.

Table 2.1. Regular-Expression Operator Precedence

operator	description	level
()	parentheses for grouping	1 (highest)
[]	character classes	2
* + ?	closure: 0 or more, 1 or more, 0 or 1	3
cc	concatenation	4
	OR	5
^ \$	anchors to beginning and end of line	6 (lowest)

2.4.3 Regular Definitions*

There is an alternate way of describing a language's token set that takes a more grammatical approach, and which is used in many language specifications. A *regular definition* builds up a language specification using a combination of regular-expression operators and production-like specifiers. For example:

<i>keyword</i>	\rightarrow	long int double while ...
<i>digit</i>	\rightarrow	0 1 2 ... 9
<i>digit_sequence</i>	\rightarrow	<i>digit</i> +
<i>sign</i>	\rightarrow	+ -
<i>exponent_part</i>	\rightarrow	e <i>sign?</i> <i>digit_sequence</i> E <i>sign?</i> <i>digit_sequence</i>
<i>floating_constant</i>	\rightarrow	<i>digit_sequence</i> . <i>digit_sequence?</i> <i>exponent_part?</i> <i>digit_sequence?</i> . <i>digit_sequence</i> <i>exponent_part?</i> <i>digit_sequence</i> <i>exponent_part</i>

Element_{opt}:

Occasionally you see an *opt* subscript used to denote an optional element, such as *digit_{opt}* rather than *digit?*. This grammatical approach to languages is discussed in greater depth in the next chapter.

2.4.4 Finite Automata*

Recognizers.

A *recognizer* program, such as a lexical analyzer, reads a string as input and outputs *yes* if the string is a sentence in a language, *no* if it isn't. A lexical analyzer has to do more than say yes or no to be useful, so an extra layer is usually added around the recognizer itself. When a certain string is recognized, the second layer performs an action associated with that string. L^EX takes an input file comprised of regular expressions and associated actions (code). It then builds a recognizer program that executes the code in the actions when a string is recognized. L^EX builds the recognizer component of the analyzer by translating regular expressions that represent the lexemes into a *finite automaton* or *finite state machine* (usually abbreviated to “state machine” or “FSM”). Strictly speaking, an FSM consists of the following:

- A finite set of *states*.
- A set of *transitions* (or moves) from one state to another. Each transition is labeled with a character from the input alphabet.
- A special *start state*.
- A set of final or *accepting states*.

Finite automata, state machines.

States.

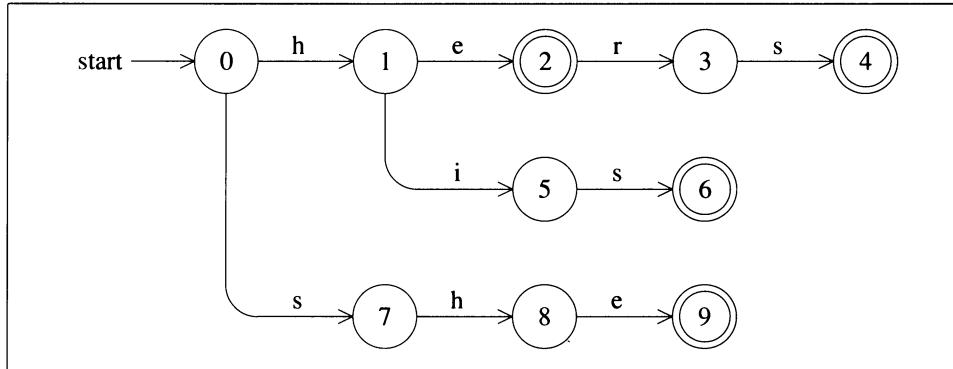
Transitions.

Start state.

Accepting states.

Transition diagram.

State machines are best understood by looking at one. Figure 2.3 is a *transition diagram* for a state machine that recognizes the four strings “he”, “she”, “his”, and “hers”.

Figure 2.3. A State Machine

The circles are individual states, marked with the state number—an arbitrary number that identifies the state. State 0 is the *start state*, and the machine is initially in this state. The lines connecting the states represent the transitions, these lines are called *edges* and the label on an edge represents characters that cause the transition from one state to another (in the direction of the arrow). From the start state, reading an *h* from the input causes a transition to State 1; from State 1, an *e* gets the machine to State 3, and an *i* causes a transition to State 5; and so on. A transition from State *N* to state *M* on the character *c* is often represented with the notation: $\text{next}(N,c)=M$. This function is called the *move* function by some authors, [Aho] among them, but I feel that *next* better describes what the function is doing.

Edges.

The states with double circles are called *accepting states*. Entering an accepting state signifies recognition of a particular input string, and there is usually some sort of action associated with the accepting state (in lexical-analyzer applications, a token is returned). Unmarked edges (for example, there are no outgoing edges marked with an *i*, *s*, *r*, or *e* from State 0) are all implied transitions to a special implicit *error state*.

next(N,c)=M.

Accepting states.

State machines such as the foregoing can be modeled with two data structures: a single variable holding the current state number and a two-dimensional array for computing the next state. One axis is indexed by the input character, the other by the current state, and the array holds the next state. For example, the previous machine can be represented by the arrays in Table 2.2. Two arrays are used, one to hold the state transitions and another to tell you whether a state is accepting or not. (You could also use a single, two-dimensional array of structures, one element of which was the next state and the other of which was the accepting-state marker, but that would waste space.) The next state is determined from the current state and input character, by looking it up in the table as follows:

Modeling state machines with arrays.

```

next_state = Transition_table[ input_character ][ current_state ];

if( Accepting[ next_state ] == 1 )
    do_an_accepting_action( next_state );
  
```

This input character is usually called the *lookahead* character because it's not removed from the input until the next-state transition is made. The machine derives the next state from the current state and lookahead character. If the next state is not the error state, then set the current state to that state and advance past the lookahead character (typically by reading, and discarding, it).

Lookahead character used to compute next state.

The machine we just looked at is called a *deterministic finite automaton* or *DFA*. A DFA is “deterministic” in that the next state can always be determined by knowing the current state and the current lookahead character. To be more specific, a DFA is a state

Deterministic finite automata (DFA).

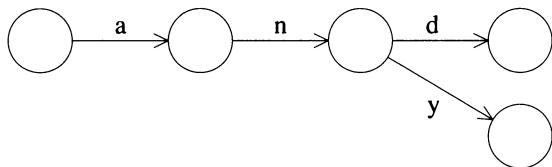
Table 2.2. Representing the State Machine

		Transition Table					Accepting	
		Lookahead Character						
		e	h	i	r	s		
current state	0	—	1	—	—	7	0	
	1	2	—	5	—	—	0	
	2	—	—	—	3	—	1	
	3	—	—	—	—	4	0	
	4	—	—	—	—	—	1	
	5	—	—	—	—	6	0	
	6	—	—	—	—	—	1	
	7	—	8	—	—	—	0	
	8	9	—	—	—	—	0	
	9	—	—	—	—	—	1	

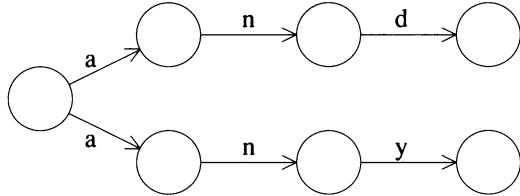
Nondeterministic finite automaton (NFA).

ϵ transitions match empty string.

machine in which all outgoing edges are labeled with an input character, and no two edges leaving a given state have the same label. There is also a second, more general type of state machine called a *nondeterministic finite automaton* or *NFA*, which is more useful in many applications, including the current one. (All DFAs are also NFAs, but not the other way around.) An NFA has no limitations on the number and type of edges: Two outgoing edges can have the same label, and edges can be labeled with the empty string, ϵ . This last type of edge is called an *epsilon edge* or *epsilon transition*. Since an ϵ transition matches an empty string, it is taken without advancing the input and is always taken—regardless of the input character. For example, how can the regular expression (*and/any*) be represented as a state machine? A DFA for this expression looks like this:

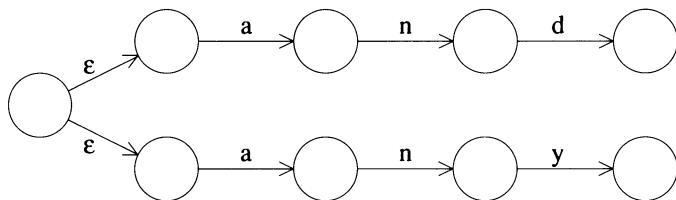


Unfortunately, DFA's are difficult to construct directly from regular expressions⁷—NFA's are easy to construct. Two possibilities are:



and

7. It is possible to construct a DFA directly from a regular expression, though I won't discuss how to do it here. See both [McNaughton] and [Aho] pp. 135–141.



The second machine is preferable because it's easier to represent in a computer program (we'll see how in a moment). As you can see, the NFA can be an awkward data structure to use. It can have many more states than the equivalent DFA, and it's difficult to write a state-machine driver (a program that uses the state machine to do something, such as recognize tokens) that can use it directly.⁸ L^EX solves the problem by creating the state machine in a two-step process. It first makes an NFA representing the input regular expressions, and it then converts that NFA to a DFA, which it in turn outputs. I'll discuss how L^EX performs this feat later in this chapter.

The state-machine representations we've been looking at are, of course, just one of many ways to represent them. You can generalize the definitions for NFAs and DFAs. A *nondeterministic finite automaton*, or NFA, is a mathematical model consisting of:

- (1) A set of states, S .
- (2) A special state in S called the *start state*. The machine is initially in this state.
- (3) A set of states in S called *accepting states*, entry into which denotes recognition of a string. Some sort of action is usually associated with each of the accepting states.
- (4) A set of input symbols (an input alphabet).
- (5) A *next* function that, when given a state and an input symbol, returns the set of states to which control is transferred on that symbol from the indicated state. I'll describe this *next* function in greater detail in a moment—note, however, that the *next* function returns a set of states. The main implementation difference between an NFA and a DFA is this *next* function. The *next* function for a DFA always yields a single next state. The equivalent NFA function can yield several next states.

A *deterministic finite automaton* or DFA is an NFA with the following restrictions:

State-machine driver.

NFA: formal definition.

- (1) No state can have an outgoing ϵ transition (an edge labeled with ϵ , the empty string).
- (2) There may be no more than one outgoing transition from any state that is labeled with the same character.

DFA: formal definition.

In practical terms, the foregoing definition describes only the data structures (the set of states and the way that the transitions are represented) and the *next* function that determines the next state from the current one. In other words, it tells us how to make a transition matrix. There is no information here about how the state machine is used; and automata can, in fact, be used in different ways depending on the application. The state machine is itself distinct from the *driver* program—the program that uses that machine.

The state machine and driver are distinct.

State-machine driver.

8. In fact, a theoretical NFA often has fewer states than an equivalent DFA because it can have more edges leaving a single state than the DFA has. Nonetheless, this sort of NFA is difficult to represent in a computer program because it has an indeterminate number of outgoing edges. The NFAs discussed in the current chapter all have more states than the equivalent DFA's because the extra states help smooth over these difficulties. I'll show how in a moment.

Using state machines for lexical analysis.

2.4.5 State-Machine-Driven Lexical Analyzers*

This section demonstrates how state machines are used for lexical analysis by looking, at a high level, at the method used by a LEX-generated lexical analyzer. I'll describe a simple table-driven lexical analyzer that recognizes decimal and floating-point constants. The following regular expressions describe these constants:

```
[0-9]+
([0-9]+|[0-9]*\.[0-9]+|[0-9]+\.\[0-9]*)(e[0-9]+)?  
return ICON;  
return FCON;
```

The code to the right of the regular expression is executed by the lexical analyzer when an input string that matches that expression is recognized. The first expression recognizes a simple sequence of one or more digits. The second expression recognizes a floating-point constant. The `(e[0-9]+)?` at the end of the second regular expression is the optional engineering notation at the end of the number. I've simplified by not allowing the usual + or - to follow the `e`, and only a lower-case `e` is recognized. The

```
([0-9]+ | [0-9]*\.[0-9]+ | [0-9]+\.\[0-9]*)
```

recognizes one of three patterns (I've added the spaces to clarify what's going on—they're not really there): The `[0-9]+` is a simple sequence of decimal digits. It's for numbers like `10e3`. Because of the way that LEX works, the `[0-9]+` on the previous line of the input specification takes precedence over the current one—an **ICON** is returned if a number does not have a trailing `e`, otherwise an **FCON** is returned. The `[0-9]*\.[0-9]+` recognizes numbers with at least one digit to the right of the decimal point, the `[0-9]+\.\[0-9]*` recognizes numbers with at least one digit to the left. You can't use `[0-9]*\.[0-9]*` because that pattern would accept a decimal point without numbers on either side. All of the following numbers are accepted:

```
1.2 1. .1 1.2e3 2e3 1
```

and, of these, the last is an **ICON** and the others are **FCONs**.

LEX uses a state-machine approach to recognize regular expressions, and a DFA that recognizes the previous expressions is shown in Figure 2.4. The same machine is represented as an array in Table 2.3. The next state is computed using that array with:

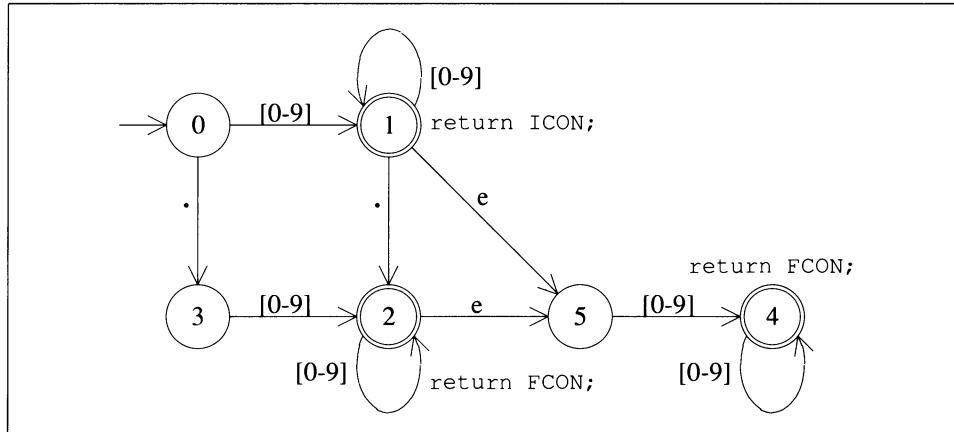
```
next_state = array[ current_state ][ input ]
```

Transition matrix.

A dash indicates a failure transition (no legal outgoing transition on the current input character from the current state). This array is typically called a *transition matrix* or *transition table*. There are three accepting states (states from which a token is recognized) in the machine: 1, 2, and 4. State 1 accepts an integer constant, and the other two recognize floating-point constants. The accepting states are recognized in an auxiliary array that is also indexed by state number, and which indicates whether or not a state is accepting.

The greedy algorithm
(matches longest string).

As I mentioned earlier, the state machine itself and the driver program that uses that machine are distinct from one another. Two algorithms are commonly used in lexical analysis applications, and the same state machine (transition matrix) is used by both algorithms. A *greedy* algorithm, shown in Table 2.4, is used by LEX (because that's what's required by most programming-language specifications). This algorithm finds the longest possible sequence of input characters that can form a token. The algorithm can be stated informally as follows: If there's an outgoing transition from the current state, take it. If the new state is an accepting state, remember it along with the input position. If there's no outgoing transition (the table has a dash in it), do the action associated with the most-recently seen accepting state. If there is no such state, then an error has occurred (LEX just ignores the partially-collected lexeme and starts over from State 0, in this situation).

Figure 2.4. State Machine That Recognizes Floating-Point Constants**Table 2.3.** State Machine in Figure 2.3 Represented as an Array

	lookahead character			accepting action
	.	0–9	e	
current state	0	3	1	—
	1	2	1	return ICON;
state	2	—	2	return FCON;
	3	—	2	—
	4	—	4	return FCON
	5	—	4	—

I'll do two examples to show the workings of the machine, the first with the input *1.2e4*. LEX starts in State 0. The 1 causes a transition to State 1 and the input is advanced. Since State 1 is a potential accepting state, the current input position and state number is remembered. The dot now gets us to State 2 (and the input is advanced again). Since State 2 is also an accepting state, the previously remembered input position and state number are overwritten by the current ones. The next input character (the 2) causes us to go from State 2 to itself. State 2 is still an accepting state, so the current input position overwrites the previously saved one. The *e* now gets us to State 5, which isn't an accepting state, so no other action is performed; and the final 4 causes a transition to State 4, which is an accepting state, so the current input position overwrites the previous one. The next input character is the end-of-input marker. There is no legal transition out of State 4 on end of input, so the machine enters the failure state. Here, the action associated with the most-recently seen accepting state (4) is performed and the machine returns FCON. The next time the subroutine is called, it returns zero immediately, because the lookahead character is end of input.

The second example looks at the incorrect input *1.2e*, with no number following the *e*. This input causes a failure transition from State 5, because there's no legal outgoing transition on end of input from State 5. When the failure occurs, the most recently seen accepting state is State 2, so the input is backed up to the condition it was in in State 2 (the next input character is an *e*) and an FCON is returned. The next time the algorithm is entered, there will be a failure transition from the start state, because an *e* can't occur at the beginning of the number. The *e* is discarded, and the algorithm goes to State 0 (and terminates).

Example: 1.2e4.

Bad-input example: 1.2e.

Table 2.4. Algorithm Used by the L^EX State-Machine Driver

```

current_state = 0;
previously_seen_accepting_state = none_seen;

if( lookahead character is end-of-input )
    return 0;

while( lookahead character is not end-of-input )
{
    if( there is a transition from the current state on the current lookahead character)
    {
        current_state = that state;
        advance the input;

        if( the current state is an accepting state )
        {
            remember the current position in the input
            and the action associated with the current state;
        }
    }
    else
    {
        if( no accepting state has been seen )
        {
            There's an error:
            Discard the current lexeme and input character.
            Current_state = 0;
        }
        else
        {
            back up the input to the position it was in when it saw the last accepting state
            perform the action associated with that accepting state;
        }
    }
}

```

Disadvantages of greedy algorithm.

Note that the greedy algorithm does have its disadvantages: It's tricky to implement and tends to be relatively slow. It can also cause the recognizer to behave in sometimes unexpected ways. (The L^EX input expression $(\backslash n | .)^*$ tries to absorb the entire input file, for example.) It is nonetheless the best (and sometimes the only) choice in most real lexical-analysis applications.

The nongreedy algorithm (matches first string).

The second type of algorithm (the *nongreedy* algorithm) is much simpler. Here, the shortest possible input string is recognized, and the machine just accepts as soon as an accepting state is entered. A nongreedy recognizer program is much simpler to implement than a greedy one, and is much faster as well. Nonetheless, this algorithm can be used only when all the accepting states in the machine are terminal nodes—when they have no outgoing transitions.

Terminal nodes.

2.4.6 Implementing a State-Machine-Driven Lexical Analyzer

This section shows how the machine in the previous section is implemented by analyzing a L^EX output file in depth. You may want to skip over this section if you’re not interested in this level of detail. You should read Appendixes A and D, in which various support functions and L^EX itself are described, before continuing. A L^EX input file that recognizes the floating-point constants we’ve been looking at is shown in Listing 2.11

L^EX input file for floating-point constants.

Listing 2.11. *numbers.lex*—A L^EX Input File to Recognize Floating-Point Constants

```

1  %{
2      #define FCON  1;
3      #define ICON  2;
4  %}
5  D      [0-9]                                /* a single decimal digit */
6  %%
7  {D}+
8  ({D}+|{D}*.{D}+|{D}+\.{D}*) (e{D}+)?    return ICON;
9  %%                                         return FCON;

```

The lex output file—*lexyy.c*—begins in Listing 2.12.⁹ The first two lines are from the header portion of the original L^EX input file. They are followed by a comment that describes the state machine that L^EX created. If a state is an accepting state, the first few characters of the equivalent code are printed (tabs are mapped to \t), along with the input line number. The goto transitions are shown along with the characters that cause the transitions. If several edges (outgoing transitions) all go to the same state, they are represented like this:

```
goto 2 on 0123456789
```

The state goes to State 2 if the input character is a digit. The entire comment is surrounded with an **#ifdef _NEVER_** in case a */ should accidentally come up in one of these lists of transition characters.

State-machine description in L^EX output file.

The next part of the L^EX output file is copied directly from the template file (*/lib/lex.par* by default, but a different file can be specified with the *-m* command-line switch). This template file is separated into three parts by formfeed (Ctrl-L) characters. Everything from the beginning of the file up to the first formfeed is copied into the output file at this juncture. The relevant code is shown in Listing 2.13.

NEVER.

The **#ifndef** directive on line 35 of Listing 2.13 lets you define YYPRIVATE in the L^EX-input-file header, without having to **#undef** it first. Most of the global variables in the file are declared as YYPRIVATE, which normally translates to the keyword **static**. Redefining this macro to an empty string makes these variables true globals, which can be accessed from outside the current file. I’m using the definition of NULL on line 39 to determine if *<stdio.h>* was included previously, and including it if not. Finally, if YYDEBUG is defined, various debugging diagnostics are activated. These are printed only if the variable yydebug is also true, thus the **if** statement on line 45. It’s best to

Template-file organization, *lex.par*

YYPRIVATE.

YY_D, yydebug,
YYDEBUG

9. Note that there are two kinds of text in *lexyy.c*: (1) text copied verbatim from the template file *lex.par* and (2) text generated by L^EX itself. The listings that describe those parts of *lexyy.c* that are copied from the template file are labeled *lex.par* in the following discussion. L^EX-generated text is in listings labeled *lexyy.c*. Line numbers carry over from one listing to another because there’s really only a single output file.

Listing 2.12. *lexyy.c*—State-Machine Description

```

1     #define FCON  1
2     #define ICON  2
3
4 #ifdef __NEVER__
5 /-----
6   * DFA (start state is 0) is:
7   *
8   * State 0 [nonaccepting]
9   *   goto 3 on .
10  *   goto 1 on 0123456789
11  * State 1 [accepting, line 7 <return ICON;>]
12  *   goto 2 on .
13  *   goto 1 on 0123456789
14  *   goto 5 on e
15  * State 2 [accepting, line 8 <return\tFCON;>]
16  *   goto 2 on 0123456789
17  *   goto 5 on e
18  * State 3 [nonaccepting]
19  *   goto 2 on 0123456789
20  * State 4 [accepting, line 8 <return\tFCON;>]
21  *   goto 4 on 0123456789
22  * State 5 [nonaccepting]
23  *   goto 4 on 0123456789
24 */
25
26 #endif

```

see how the macro works by looking at an example; if YYDEBUG is defined, then a debugging diagnostic like:

```
YY_D( printf("aaaaagh!!!") );
```

is expanded to:

```
if( yydebug ){ printf("aaaaagh!!!"); } else;
```

Trailing **else** in
multiple-statement
macro.

Note that the semicolon following the **else** comes from the original macro invocation and the semicolon following the `printf()` follows the `x` in the macro definition. That trailing **else** is important in order to make something like the following work correctly:

```
if( something )
    YY_D(printf("aaaaagh!!!") );
else
    something_else();
```

The foregoing expands to:

```
if( something )
    if( yydebug )
    {
        printf("aaaaagh!!!");
    }
    else
    ;
else
    something_else();
```

If the **else** weren't present in the macro definition, then the **else** `something_else()` clause in the original code would incorrectly bind to the

Listing 2.13. *lex.par*— Various Definitions Copied Into *lexyy.c*

```

27 /* YY_TTYPE is used for the DFA transition table: Yy_nxt[], declared below.
28 * YYF marks failure transitions in the DFA transition table. There's no failure
29 * state in the table itself, these transitions must be handled by the driver
30 * program. The DFA start state is State 0. YYPRIVATE is only defined here only
31 * if it hasn't be #defined earlier. I'm assuming that if NULL is undefined,
32 * <stdio.h> hasn't been included.
33 */
34
35 #ifndef YYPRIVATE
36 #define YYPRIVATE static
37#endif
38
39 #ifndef NULL
40 #include <stdio.h>
41#endif
42
43 #ifdef YYDEBUG
44     int yydebug = 0;
45 #define YY_D(x) if( yydebug ){ x; }else
46 #else
47 #define YY_D(x)
48#endif
49
50 typedef unsigned char YY_TTYPE;
51 #define YYF ((YY_TTYPE)(-1))
52
53 unsigned char *ii_text();

```

`if(yydebug)` rather than the `if(something)`. If `YYDEBUG` isn't defined in the header, then the argument to `YY_D` effectively disappears from the input (the macro expands to an empty string). In this case, the `printf()` statements go away.

The code on lines 50 and 51 of Listing 2.13 are used to declare and access the transition-matrix array. `YY_TYPE` is the type of one array element, and `YY_F` marks failure transitions in the array. This latter value cannot be used as a state number. Note that the cast to `unsigned char` effectively translates `-1` to 255. Similarly, `-1`'s in the tables are all silently converted to 255 as part of the initialization.

The next part of the LEX output file is the state-machine transition matrix. It is used to compute the next state from the current state and lookahead symbols. This array can take three forms. The first, uncompressed form is shown in Figure 2.5 and Listing 2.14. I've simplified the picture by leaving all the error transitions blank. (They're initialized to `-1` in Listing 2.14.) An uncompressed array is generated by specifying a `-f` (for fast) switch on the LEX command line. The next state is computed with:

```
Yy_nxt[ current_state ][ lookahead_character ]
```

This operation is encapsulated into the `yy_next(state,c)` macro on line 147 of Listing 2.14.

Notice that several columns in the uncompressed array are identical. All columns not associated with a period, `e`, or digit are the same—they're all error transitions. All the columns for the digits are the same. By the same token, the rows associated with States 4 and 5 are the same. (The states aren't equivalent because one is an accepting state and the other isn't.) This situation holds with most state machines that recognize real token sets. Taking C as a case in point, all the control characters and the space character are ignored, so the columns for these are identical. With a few exceptions like `L`

`YY_TYPE`, `YYF`.

Transition matrix representations.

Uncompressed transition matrix.

Compressed transition matrix. Redundant row and column elimination.

Figure 2.5. The Uncompressed Transition Table

Yy_nxt[][]	.	0123456789	e	
0	3	1111111111		
1	2	1111111111	5	return ICON
2		2222222222	5	return FCON
3		2222222222		
4		4444444444		
5		4444444444		return FCON

Listing 2.14. lexyy.c — The Uncompressed Transition Table

```

54 YYPRIVATE YY_TTYPE Yy_nxt[ 6 ][ 128 ] =
55 {
56 /* 00 */ { -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
57 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
58 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
59 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
60 -1, -1, -1, -1, -1, -1, -1, 3, -1, 1, 1, 1,
61 1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -1, -1,
62 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
63 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
64 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
65 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
66 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
67 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
68 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1
69 },
70 /* 01 */ { -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
71 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
72 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
73 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
74 -1, -1, -1, -1, -1, -1, -1, 2, -1, 1, 1, 1,
75 1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -1, -1,
76 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
77 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
78 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
79 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
80 -1, 5, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
81 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
82 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1
83 },
84 /* 02 */ { -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
85 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
86 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
87 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
88 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, 2, 2,
89 2, 2, 2, 2, 2, 2, 2, 2, -1, -1, -1, -1,
90 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
91 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
92 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
93 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
94 -1, 5, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
95 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
96 -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1

```



Listing 2.14. continued...

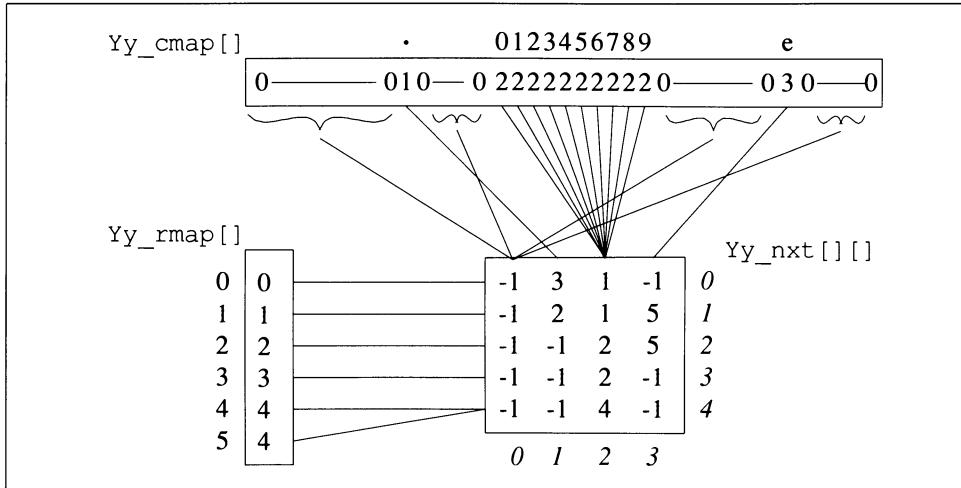
```

97      },
98  /* 03 */ { -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
99    -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
100   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
101   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
102   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, 2, 2,
103   2, 2, 2, 2, 2, 2, 2, 2, -1, -1,
104   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
105   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
106   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
107   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
108   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
109   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
110   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1
111 },
112 /* 04 */ { -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
113   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
114   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
115   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
116   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, 4, 4,
117   4, 4, 4, 4, 4, 4, 4, 4, -1, -1,
118   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
119   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
120   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
121   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
122   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
123   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
124   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1
125 },
126 /* 05 */ { -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
127   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
128   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
129   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
130   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, 4, 4,
131   4, 4, 4, 4, 4, 4, 4, 4, -1, -1,
132   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
133   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
134   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
135   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
136   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
137   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
138   -1, -1, -1, -1, -1, -1, -1, -1, -1, -1
139 }
140 };
141
142 /*****
143 * yy_next(state,c) is given the current state and input
144 * character and evaluates to the next state.
145 */
146
147 #define yy_next(state, c) Yy_nxt[ state ][ c ]

```

and x , all the columns for the letters are identical, as are most of the digit's columns. Moreover, at least half of the states in a typical machine have no legal outgoing transitions, so the rows associated with these states are identical—every cell holds -1 .

LEX's default compression technique takes advantage of this situation and eliminates the redundant rows and columns by creating two supplemental arrays. The compressed

Figure 2.6. Transition Table With Redundant Rows and Columns Eliminated

`Yy_cmap[]`, `Yy_nxt[]`

`Yy_rmap[]`.

table is shown in Figure 2.6 and in Listing 2.15. The `Yy_cmap[]` array is indexed by lookahead character and holds the index of one of the columns in the `Yy_nxt[]` array. When several columns in the original array are equivalent, the matching entries in `Yy_cmap[]` hold the index of a single column in `Yy_nxt[]`. For example, the columns associated with digits in the original table are all identical. Only one of these columns is present in the compressed array (at `Yy_nxt[x][2]`), and all columns corresponding to digits in `Yy_cmap` hold a 2. The rows are compressed in the same way using `Yy_rmap[]`. Since rows 4 and 5 are identical in the uncompressed array, `Yy_rmap[4]` and `Yy_rmap[5]` both hold a 4, and `Yy_nxt[4][x]` holds the original row from the uncompressed table.

An array element is accessed using:

```
Yy_nxt[ Yy_rmap[current_state] ][ Yy_cmap[lookahead_character] ]
```

rather than the

```
Yy_nxt[ current_state ][ lookahead_character ]
```

that's used for the uncompressed table. The `yy_next` macro for this type of table is defined on line 102 of Listing 2.15.

Compression ratio.

Redundant-row-and-column elimination is usually the best practical compression technique. The access time is fast, and the compression ratio is usually quite good. (In this example, the ratio is about 4:1 — 154 bytes, as compared to 640 bytes. The C lexical analyzer presented in Appendix D does even better, with a compression ratio of about 7:1 — 1,514 bytes versus 10,062 bytes.)

Pair-compressed transition matrix.

`Yy_nxtN`, `Yy_nxt[]`.

A second compression method yields better compression if the transition matrix is particularly sparse, though the access time is slower. The rows are split into distinct one-dimensional arrays, accessed indirectly through an array of pointers (see Figure 2.7). The rows are all named `Yy_nxtN`, where `N` is the original row index in the uncompressed table (row 5 is in `Yy_nxt5[]`), and the array of pointers is called `Yy_nxt[]`. The current tables are compressed in this way in Listing 2.16.

If the first cell of the `Yy_nxtN` array is zero, then the remainder of the array is identical to the original row, and can be accessed directly using the lookahead character. To simplify a little:

Listing 2.15. *lexyy.c*— Transition Table With Redundant Rows and Columns Eliminated

```

54  /-----
55  * The Yy_cmap[] and Yy_rmap arrays are used as follows:
56  *
57  * next_state= Yy_nxt[ Yy_rmap[current_state] ][ Yy_cmap[input_char] ];
58  *
59  * Character positions in the Yy_cmap array are:
60  *
61  *   ^@  ^A  ^B  ^C  ^D  ^E  ^F  ^G  ^H  ^I  ^J  ^K  ^L  ^M  ^N  ^O
62  *   ^P  ^Q  ^R  ^S  ^T  ^U  ^V  ^W  ^X  ^Y  ^Z  ^[  ^\  ^]  ^^  ^
63  *   !   "   #   $   %   &   '   (   )   *   +   ,   -   .   /
64  *   0   1   2   3   4   5   6   7   8   9   :   ;   <   =   >   ?
65  *   @   A   B   C   D   E   F   G   H   I   J   K   L   M   N   O
66  *   P   Q   R   S   T   U   V   W   X   Y   Z   [   \   ]   ^   -
67  *   `   a   b   c   d   e   f   g   h   i   j   k   l   m   n   o
68  *   p   q   r   s   t   u   v   w   x   y   z   {   |   }   ~   DEL
69 */
70
71 YYPRIVATE YY_TTYPE Yy_cmap[128] =
72 {
73     0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,
74     0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,
75     0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0,
76     2, 2, 2, 2, 2, 2, 2, 2, 2, 2, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,
77     0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,
78     0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,
79     0, 0, 0, 0, 0, 3, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0,
80     0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0
81 };
82
83 YYPRIVATE YY_TTYPE Yy_rmap[6] =
84 {
85     0, 1, 2, 3, 4, 4
86 };
87
88 YYPRIVATE YY_TTYPE Yy_nxt[ 5 ][ 4 ]=
89 {
90 /* 00 */ { -1, 3, 1, -1 },
91 /* 01 */ { -1, 2, 1, 5 },
92 /* 02 */ { -1, -1, 2, 5 },
93 /* 03 */ { -1, -1, 2, -1 },
94 /* 04 */ { -1, -1, 4, -1 }
95 };
96
97 /-----
98 * yy_next(state,c) is given the current state number and input
99 * character and evaluates to the next state.
100 */
101
102 #define yy_next(state,c) (Yy_nxt[ Yy_rmap[state] ][ Yy_cmap[c] ])

```

```

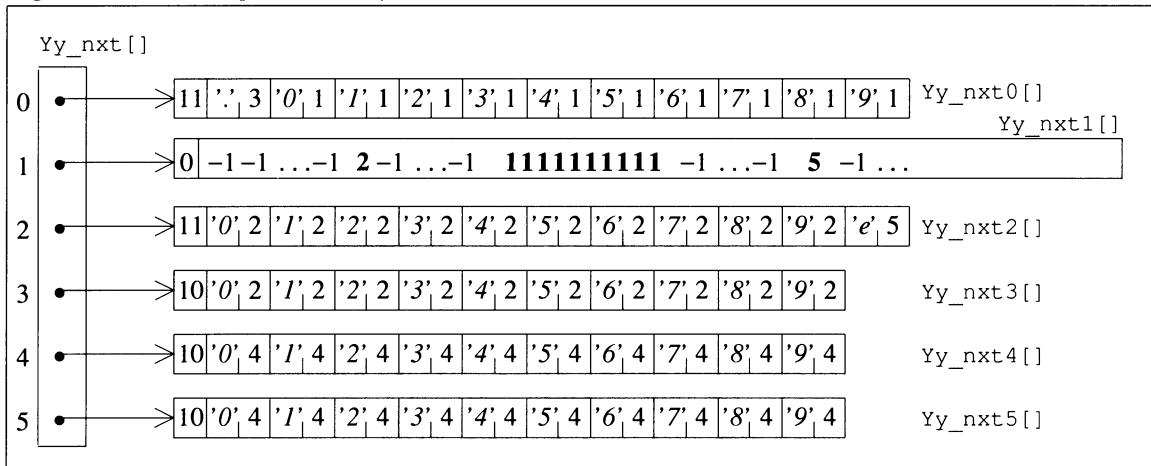
YY_TTYPE *row;
row = Yy_nxt[ current_state ];

if( *row == 0 )
    next_state = (row + 1)[ lookahead_character ];

```

If the first cell of the *Yy_nxtN* array is nonzero, then the array holds a sequence of character/next-state pairs and the first cell holds the number of pairs. For example,

Figure 2.7. Pair-Compressed Transition Table



Listing 2.16. *lexyy.c*—Pair-Compressed Transition Table

```

44 YYPRIVATE YY_TTYPE Yy_nxt0 [] = { 11,
45     '.', 3, '0', 1, '1', 1, '2', 1, '3', 1,
46     '4', 1, '5', 1, '6', 1, '7', 1, '8', 1,
47     '9', 1};
48 YYPRIVATE YY_TTYPE Yy_nxt1 [] = { 0,
49     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
50     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
51     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
52     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
53     -1, -1, -1, -1, -1, -1, -1, -1, 2, -1, 1, 1,
54     1, 1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -1,
55     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
56     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
57     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
58     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
59     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
60     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
61     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
62     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
63     -1, -1, -1, -1, -1, -1, -1, -1, 2, -1, 1, 1,
64     1, 1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -1,
65     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
66     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
67     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
68     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
69     -1, 5, -1, -1, -1, -1, -1, -1, -1, -1, -1,
70     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1,
71     -1, -1, -1, -1, -1, -1, -1, -1, -1, -1, -1 };
72 YYPRIVATE YY_TTYPE Yy_nxt2 [] = { 11,
73     '0', 2, '1', 2, '2', 2, '3', 2, '4', 2,
74     '5', 2, '6', 2, '7', 2, '8', 2, '9', 2,
75     'e', 5};
76 YYPRIVATE YY_TTYPE Yy_nxt3 [] = { 10,
77     '0', 2, '1', 2, '2', 2, '3', 2, '4', 2,
78     '5', 2, '6', 2, '7', 2, '8', 2, '9', 2
79   };
80 YYPRIVATE YY_TTYPE Yy_nxt4 [] = { 10,
81     '0', 4, '1', 4, '2', 4, '3', 4, '4', 4,
82     '5', 4, '6', 4, '7', 4, '8', 4, '9', 4
83   };
84 YYPRIVATE YY_TTYPE Yy_nxt5 [] = { 10,
85     '0', 4, '1', 4, '2', 4, '3', 4, '4', 4,
86     '5', 4, '6', 4, '7', 4, '8', 4, '9', 4
87   };
88
89 YYPRIVATE YY_TTYPE *Yy_nxt[ 6 ] =
90 {
91     Yy_nxt0 , Yy_nxt1 , Yy_nxt2 , Yy_nxt3 , Yy_nxt4 , Yy_nxt5

```

Listing 2.16. continued...

```

92  };
93
94 /*-----*/
95
96 YYPRIVATE YY_TTYPE yy_next( cur_state, c )
97 unsigned int c ;
98 int cur_state ;
99 {
100    /* Given the current state and the current input character
101       * return the next state.
102       */
103
104    YY_TTYPE *p = Yy_nxt[ cur_state ] ;
105    register int i;
106
107    if( p )                                /* there are transitions */
108    {
109        if( (i = *p++) == 0 )                /* row is uncompressed */
110            return p[ c ];
111
112        for( ; --i >= 0 ; p += 2 )          /* row is in pairs */
113            if( c == p[0] )
114                return p[1];
115    }
116    return YYF;
117 }
```

`Yy_nxt0[]` in Figure 2.7 contains 11 pairs (because 11 is in the first cell). The first pair is `{'.',3}`, which means that if the next input character is a dot, then the next state is State 3. The next pair is `{'0',1}`—if the lookahead character is a '0', go to State 1, and so forth. If you get to the end of the list without finding the current lookahead character, there is a failure transition on that character. Again, simplifying a little:

```

YY_TTYPE *row;
row = Yy_nxt[ current_state ];

if( *row != 0 )
{
    for( num_pairs = *row++ ; --num_pairs >= 0 ; row += 2)
        if( row[0] == lookahead_character )
        {
            next_state = row[1];
            break;
        }
}
```

The foregoing is implemented by the code on line 112 of Listing 2.16.

A third situation, unfortunately not illustrated in this example, occurs when a state has no legal outgoing transitions. In this case, `Yy_nxt[state]` is set to NULL, so:

```

if( Yy_nxt[current_state] == NULL )
    next_state = FAILURE;
```

This transition is activated when the test on line 107 of Listing 2.16 fails.

Note that further compression can be achieved at the expense of error recovery by providing a default state other than the error state. Hitherto, an error was indicated when you got to the end of the pair list. You can use the most common nonerror transition instead of the error transition in this case, however. For example, `Yy_nxt2` contains

Using a default transition.

these pairs:

[‘0’,2]	[‘1’,2]	[‘2’,2]	[‘3’,2]	[‘4’,2]
[‘5’,2]	[‘6’,2]	[‘7’,2]	[‘8’,2]	[‘9’,2]
				[‘e’,5]

All but one of the transitions are to State 2, so if you use the transition to State 2 as the default (rather than a transition to the error state), the row could be compressed as follows

```
YYPRIVATE YY_TTYPE Yy_nxt2 [] = { 1, 2, 'e', 5 };
```

The first number (1) is the pair count, as before. The second number (2) is the default next-state transition. If the *next* function gets through the entire list of pairs without finding a match, it will go to this default next state. The remainder of the array is the one pair that doesn’t go to the default state. This extreme compression isn’t very useful in a lexical-analyzer application because you can’t really afford to discard all the error information, but it’s sometimes useful in parser applications. The UNIX *yacc* utility uses a variation on this method to compress its tables.¹⁰ Of course, if there are more transitions to an explicit state than to the error state, you can put the error transitions into the pair list and use the explicit state as the default state. This way you won’t lose any information.

Setting pair-compression threshold in LEX.

Pair compression is activated in LEX with a *-cN* command-line switch. *N* is the threshold beyond which pairs are abandoned in favor of a simple array indexed by lookahead character. The example we’ve been looking at had the threshold set at 11, so any state with more than 11 outgoing transitions is handled with a simple array, and states with 11 or fewer outgoing transitions are represented with character/next-state pairs. The default threshold—used when no *N* is given on the command line—is four.

Pair compression, compression ratio.

The compression ratio here tends not to be as good as with redundant-row-and-column elimination in programming-language applications. The current example uses 247 bytes, versus 154 for the other method. The C lexical analyzer uses 3,272 bytes for the pair-compressed tables, versus 1,514 for the default method. It does do better when the data in the matrix is both sparse and randomly distributed, however. The *-v* command-line switch to LEX causes the final table sizes to be printed, so you can judge which method is more appropriate for a given application.

Accepting-state array, *Yyaccept* [].

You’ll note that the redundant-row-and-column elimination could be combined with the pair-compression technique. For example, since the last two rows in the table shown in Figure 2.7 are the same, you really need to store only one of them and keep two pointers to it. I haven’t implemented this combined method.

The next part of *lexyy.c* is the accepting-state array, *Yyaccept* [], shown in Listing 2.17. The array is indexed by state number. It evaluates to 0 if the state is not an accepting state. Other values set the conditions under which the lexeme is accepted. It holds 1 if the string is anchored to the start of the line (a ^ was the leftmost symbol in the original regular expression—an extra newline will be at the far left of the lexeme in this case). It holds 2 if the string is anchored to the end of the line (a \$ was the rightmost symbol in the original regular expression—an extra newline will be at the far right of the lexeme in this case). It holds 3 if the lexeme is anchored both to the beginning and the end of the line, and 4 if the lexeme is always accepted—no anchors were present in the original regular expression.

10. The method is described in [Aho], pp. 144–146. They use separate arrays for the pairs and the default transitions, but the rationale is the same.

Listing 2.17. *lexyy.c*—Accepting-State Identification Array

```

118  /*-----*
119  * The Yyaccept array has two purposes. If Yyaccept[i] is 0 then state
120  * i is nonaccepting. If it's nonzero then the number determines whether
121  * the string is anchored, 1=anchored at start of line, 2=at end of
122  * line, 3=both, 4=line not anchored
123  */
124
125 YYPRIVATE YY_TTYPE Yyaccept[] =
126 {
127     0 , /* State 0 */
128     4 , /* State 1 */
129     4 , /* State 2 */
130     0 , /* State 3 */
131     4 , /* State 4 */
132     0 , /* State 5 */
133 };

```

The remainder of *lexyy.c* file is the actual state-machine driver, shown in Listing 2.18. The first and last part of this listing are the second and third parts of the Ctrl-L-delimited template file discussed earlier. The case statements in the middle (on lines 287 to 295 of Listing 2.18) correspond to the original code attached to the regular expressions in the input file and are generated by L^EX itself.

The various global variables that communicate with the parser are declared on lines 138 to 141. Note that *yyout* is provided for UNIX compatibility, but you shouldn't use it if you're using *occs* (because it will mess up the windows in the debugging system). Same goes for the *output()* and *ECHO* macros on lines 147 and 148. UNIX supports them but they shouldn't be used in an *occs* environment. It's best to use the actual output functions, or to supply similarly-named replacement functions that you can use to debug your lexical analyzer (assuming that the functions you supply will eventually be replaced by the *occs* versions). These replacement functions are shown in *lex_io.c* (Listing 2.18). Link this file to the lexical analyzer when you're debugging a *lex* output file without a parser. Use the versions of the routines that are in *l.lib*, and which support the *occs* debugging environment, when you're using an *occs*-generated parser.

The *YYERROR()* macro on line 151 of Listing 2.19 prints internal error messages. There's no UNIX equivalent because *lex* doesn't print error messages. In the *occs* environment, you should redefine *YYERROR()* to use *yyerror()* rather than *fprintf()*. (Do it in a '%{}%' block in the definitions section.)

The *ymore()* macro on line 154 of Listing 2.19 just sets a flag to true. It forces the driver to continue processing the current lexeme, ignoring the current accepting action. *unput()* and *yyless()* (on lines 156 and 157) are the two pushback functions. They unterminate the current lexeme, push back any requested characters, and then reterminate the lexeme. I've made extensive use of the comma operator here in order to squeeze several instructions into a single macro. The comma operator just executes the comma-delimited statements in sequence. It evaluates to the rightmost statement in the list, but neither of these macros take advantage of this behavior—they'd be declared **void** if they were functions. Braces can't be used here because of the binding problems discussed earlier. Note that the conditional on line 158 sets '*yylen*' to zero if the *ii_pushback()* call fails, as it will if you try to push back too many characters; otherwise, '*yylen*' is just reduced by the number of pushed-back characters.

The *input()* function on line 162 of Listing 2.19 is complicated enough to be a subroutine. It has to return a value, and the contortions necessary to do this with a

L^EX state-machine driver.

yyout, *output()*,
ECHO.

YYERROR().

ymore(), *yyless()*,
unput().

input().

Listing 2.18. *lex_io.c*— Debugging Output Routines for L^EX-Generated Analyzer

```

1 #include <stdio.h>
2 #include <stdarg.h>
3
4 /* This file contains two output routines that replace the ones in yydebug.c,
5 * found in l.lib and used by occs for output. Link this file to a LEX-
6 * generated lexical analyzer when an occs-generated parser is not present.
7 * Then use yycomment() for messages to stdout, yyerror() for messages to
8 * stderr.
9 */
10
11 PUBLIC void yycomment( fmt, ... )
12 char *fmt;
13 {
14     /* Works like printf(). */
15
16     va_list args;
17     va_start( args, fmt );
18     vfprintf( stdout, fmt, args );
19     va_end ( args );
20 }
21
22 PUBLIC void yyerror( fmt, ... )
23 char *fmt;
24 {
25     /* Works like printf() but prints an error message along with the
26     * current line number and lexeme.
27     */
28
29     va_list args;
30     va_start( args, fmt );
31
32     fprintf ( stderr, "ERROR on line %d, near <%s>\n", yylineno, yytext );
33     vfprintf( stderr, fmt, args );
34     va_end ( args );
35 }
```

comma operator are not worth the effort.

`yylex()`, `yystate`.

The lexical analyzer itself, `yylex()` starts on line 177 of Listing 2.19. The current state is held in `yystate`, which initially set to `-1`—a value that's not used as a state number. The code on lines 187 to 192 is executed only once because `yystate` is set to `-1` only the first time the subroutine is called. The

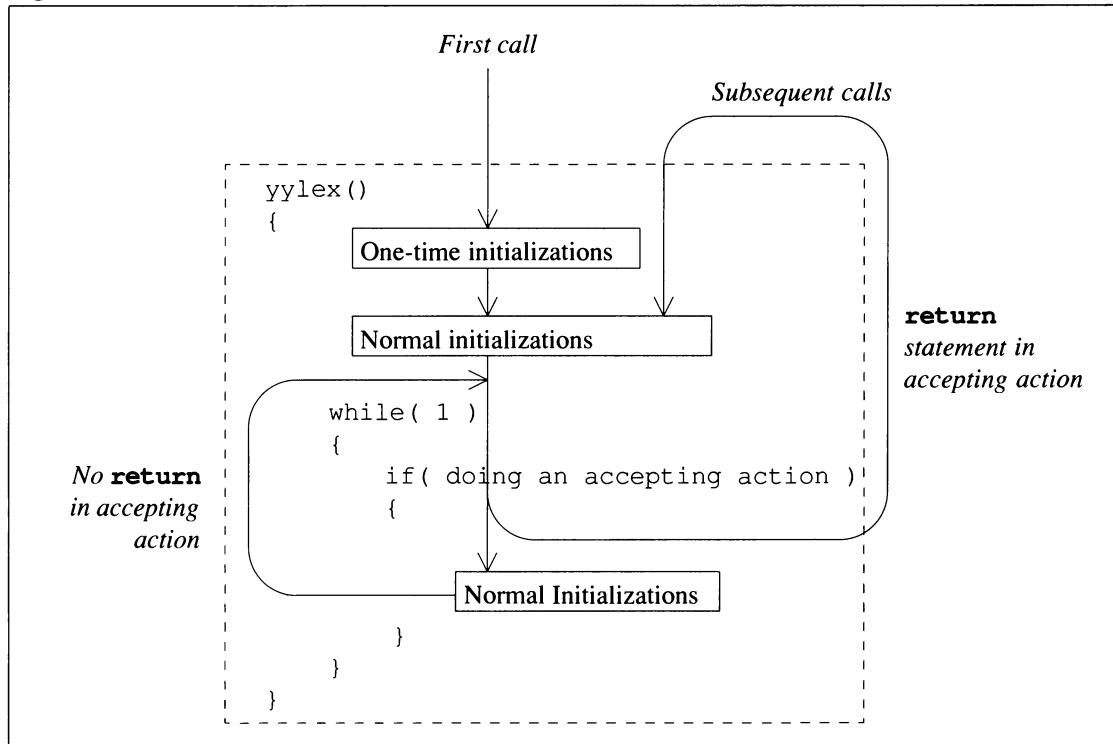
```

    ii_advance();
    ii_pushback(1);
```

forces an initial buffer load so that `ii_look()` can be used later on.

Control-flow in `yylex()`.

The actual control flow through the program is unusual in that one branch of the main loop exits the subroutine entirely and reenters the loop from the top. In other words, if an action in the original input file contains a `return` statement, then control passes out of the loop at that point, and passes back into the loop on the next call. A normal path through the loop is also available when no such `return` is executed. The situation is illustrated in Figure 2.8. The initializations can't be done at the top of the loop because they're performed only on accepting a lexeme, not on every iteration of the loop.

Figure 2.8. Flow of Control Within `yylex()`

The initializations on lines 194 to 198 of Listing 2.19 are executed every time the loop is entered from the top, and these initializations are duplicated on lines 303 to 315 for those situations where an accepting action does not contain a `return` statement. This code unterminates the previous lexeme (`ii_unterm()`), sets the start marker for the current lexeme (`ii_mark_start()`), and sets `yylastaccept` to zero to signify that no accepting state has been seen (the start state, which is always State 0, cannot be an accepting state). The `ii_unterm()` call does nothing if the string is not '\0' terminated, as is initially the case.

Listing 2.19. *lex.par*—State-Machine Driver Copied to *lexyy.c*

```

134  /*-----
135  * Global variables used by the parser.
136  */
137
138  char  *yytext;          /* Pointer to lexeme.          */
139  int   yylen;           /* Length of lexeme.          */
140  int   yylineno;         /* Input line number.          */
141  FILE  *yyout = stdout;
142
143  /*-----
144  * Macros that duplicate functions in UNIX lex:
145  */
146
147  #define output(c)  putc(c,yyout)
148  #define ECHO      fprintf(yyout, "%s", yytext )
149
  
```



Listing 2.19. continued...

```

150 #ifndef YYERROR
151 # define YYERROR(t) fprintf(stderr,"%s", t )
152 #endif
153
154 #define yymore()      yymoreflg = 1
155
156 #define unput(c)    (ii_unput(c), --yyleng )
157 #define yyless(n)   ( ii_unterm(), \
158                  ( yyleng == ii_pushback(n) ? n : yyleng ), \
159                  ii_term() \
160                )
161
162 int input()                                /* This is a macro in UNIX lex */
163 {
164     int c;
165
166     if( (c = ii_input()) && (c != -1) )
167     {
168         yytext = ii_text();
169         yylineno = ii_lineno();
170         ++yyleng;
171     }
172     return c;
173 }
174
175 /*****  

176
177 yylex()
178 {
179     int yymoreflg;           /* Set when yymore() is executed */
180     static int yystate = -1; /* Current state. */
181     int yylastaccept;       /* Most recently seen accept state */
182     int yyprev;             /* State before yylastaccept */
183     int ynstate;            /* Next state, given lookahead. */
184     int yylook;              /* Lookahead character */
185     int yyanchor;           /* Anchor point for most recently seen */
186                             /* accepting state. */
187     if( yystate == -1 )
188     {
189         yy_init_lex();        /* One-time initializations */
190         ii_advance();
191         ii_pushback(1);
192     }
193
194     yystate = 0;             /* Top-of-loop initializations */
195     yylastaccept = 0;
196     yymoreflg = 0;
197     ii_unterm();
198     ii_mark_start();
199
200     while( 1 )
201     {
202         /* Check end of file. If there's an unprocessed accepting state,
203          * yylastaccept will be nonzero. In this case, ignore EOF for now so
204          * that you can do the accepting action; otherwise, try to open another
205          * file and return if you can't.
206
207

```



Listing 2.19. continued...

```

208         while( 1 )
209     {
210         if( (yylook=ii_look(1)) != EOF )
211     {
212         yynstate = yy_next( yystate, yylook );
213         break;
214     }
215     else
216     {
217         if( yylastaccept )           /* still something to do */
218         {
219             yynstate = YYF;
220             break;
221         }
222         else if( yywrap() )        /* another file? */
223         {                           /* no */
224             yytext = "";
225             yyleng = 0;
226             return 0;
227         }
228     else
229     {
230         ii_advance();            /* load a new buffer */
231         ii_pushback(1);
232     }
233 }
234 }
235
236
237 if( yynstate != YYF )
238 {
239     YY_D( printf(" Transition from state %d", yystate) );
240     YY_D( printf(" to state %d on <%c>\n", yynstate, yylook) );
241
242     if( ii_advance() < 0 )          /* Buffer full */
243     {
244         YYERROR( "Lexeme too long, discarding characters\n" );
245         ii_flush(1);
246     }
247
248     if(yyanchor = Yyaccept[ yynstate ])    /* saw an accept state */
249     {
250         yyprev      = yystate ;
251         yylastaccept = yynstate ;
252         ii_mark_end();           /* Mark input at current character. */
253                                         /* A subsequent ii_to_mark() */
254                                         /* returns us to this position. */
255     }
256
257     yystate = yynstate;
258 }
259 else
260 {
261     if( !yylastaccept )           /* illegal input */
262     {
263 #ifdef YYBADINP
264         YYERROR( "Ignoring bad input\n" );
265 #endif
266         ii_advance();           /* Skip char that caused failure. */
267     }

```



Listing 2.19. continued...

```

268         else
269     {
270         ii_to_mark();           /* Back up to previous accept state */
271         if( yyanchor & 2 )    /* If end anchor is active          */
272             ii_pushback(1);    /* push back the CR or LF          */
273
274         if( yyanchor & 1 )      /* if start anchor is active */
275             ii_move_start();    /* skip the leading newline */
276
277         ii_term();            /* Null-terminate the string */
278         yylen = ii_length();
279         yytext = ii_text();
280         yylineno = ii_lineno();
281
282         YY_D( printf("Accepting state %d, ", yylastaccept) );
283         YY_D( printf("line %d: <%s>\n", yylineno, yytext ) );
284
285         switch( yylastaccept )
286     {
287         case 1:                  /* State 1   */
288             return ICON;
289             break;
290         case 2:                  /* State 2   */
291             return FCON;
292             break;
293         case 4:                  /* State 4   */
294             return FCON;
295             break;
296
297         default:
298             YYERROR( "INTERNAL ERROR, yylex\n" );
299             break;
300     }
301 }
302
303     ii_unterm();
304     yylastaccept = 0;
305
306     if( !yymoreflg )
307     {
308         yystate = 0;
309         ii_mark_start();
310     }
311     else
312     {
313         yystate = yyprev;      /* Back up */
314         yymoreflg = 0;
315     }
316 }
317 }
318 }
```

yymoreflg.

The yymoreflg that is tested on line 306 of Listing 2.19 is set true by yymore(). If yymoreflg is false, the machine behaves normally: the next state is State 0 and the start-of-lexeme marker is reset so that the machine can collect a new lexeme (on lines 308 and 309). If yymoreflg is true, the machine backs up one state rather than going to State 0, and the start-of-lexeme marker isn't modified, so additional characters are added

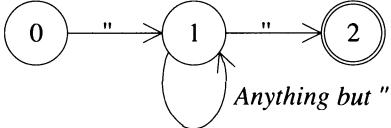
to the end of the current lexeme. The only use of `yyprev` is to remember the last state so that the machine can back up.

Generally, backing up is the correct action to take for `yymore()`. For example, the naive string-processing algorithm discussed in Chapter Two looked like this:

```
\\"[^\\"]*\\"    if( yytext[yylen-2] == '\\\\' )
                  yymore();
else
    return STRING;
```

`yyprev`Problems with
`yymore()`.

This expression creates the following machine:



The problem is that the machine is in State 2 after the close quote is recognized. In order to continue processing the string, it needs to back up to State 1 (so that another close quote can get it back to State 2). The back up that's initiated by `yymore()` can cause problems (not very often, but, as Shiva says, to be four-armed is to be forewarned). One of my original attempts to handle the escape-sequence-in-a-string problem looked like this:

```
\\"[^\\\"]*\\"    if( ii_lookahead(1) == '\\\\' )
{
    input(); /* Skip the backslash */
    input(); /* and the character that follows. */
    yymore();
}
else
{
    input(); /* Get the close quote */
    return STRING;
}
```

The idea was to break out of the regular expression if either a backslash or quote was encountered, look ahead to see which of the two possible characters were there, and then absorb the escape sequence with `input()` calls. The state machine looks like this:



The problem arises when you try to handle a string like:

"\\\""

The machine starts out in State 0, the quote puts it into State 1, and then the machine terminates because there's no outgoing transition from State 1 on a backslash. The code is now activated and the `if` statement tests true. The two `input()` calls absorb the backslash and the second quote, and `yymore()` backs us up to the previous state—State 0. Now the third quote is encountered, but the machine treats it as a start-of-string character, not as an end-of-string character. The code associated with State 1 won't be executed until a fourth quote or another backslash is encountered. The best solution to this problem prevents the backup by using:

```
yyprev = yystate;
yymore();
```

rather than a simple `yymore()` invocation. To see what's happening here, consider that `yyprev` holds the number of the state to back up to—the previous state; `yystate` holds

the current state number. The assignment of the current state to the previous state means that the machine will back up to the current state when `yymore()` is invoked. That is, it won't back up at all.

Alternately, you could use the following:

```
ii_unterm();
continue;
```

instead of `yymore`, but I think that this latter solution is more confusing. (Look at the code to see what a `continue` will do here.) You don't want to `break` out of the loop because the existing lexeme is discarded in this case.

Returning to Listing 2.19, the `while` loop on lines 208 to 234 (page 77) gets the next input character and puts it into `yylook`—the input is not advanced by `ii_look()`. The normal situation of not being at end of file is handled on line 212, where the next state is computed. The `else` clause deals with end of file. If `yylastaccept` is true on line 217, then the machine hasn't executed the accepting action for the last lexeme in the file, so end-of-file processing is delayed until the action is done. If it's false, then `yywrap()` is called to open a new input file. This user-supplied subroutine can be used to chain together several input files. The default, library version just returns zero. If you have several input files, you can replace the default version with one that opens the next input file (and returns 1 until there are no more files to open). So, if `yywrap()` returns false, it's not time to wrap up and the code loops back up to line 208. The next state is then recomputed using the first character in the new file (as if the EOF had not been encountered). The code loops until a nonempty file is opened. There's some potential confusion here in that `yywrap()` returns true if the program should terminate, even though a false return value would be more intuitive in this situation. Remember, the name `yywrap()` stands for “go ahead and wrap up.”¹¹

When the loop terminates, `yylook` holds the current lookahead character, and the potential next state is figured on line 212. The machine has not changed state yet because `yystate` hasn't been modified. The machine is looking ahead at what the next state is going to be, given the current lookahead character.

If the next state is not a failure transition, the input is advanced (on line 242) and the machine looks to see if the new state is an accepting state (on line 248). If so, the accepting state is remembered in `yylastaccept`, and the state preceding the accepting state is also remembered for `yymore()` processing. Finally, the driver switches to a new state by modifying `yystate` on line 257.

The `else` clause that starts on line 259 handles failure transitions. In this case you want to perform the accepting action associated with the most recently seen accepting state (which you just remembered in `yylastaccept`). If `yylastaccept` is zero, then no such accepting state was encountered and you're looking at a bad lexeme (one that is not described by any regular expression in the input file). An error message is printed if `YYBADINP` is true and `ii_advance()` is called to skip the offending character.

If an accepting state had been encountered, the input is restored to the condition it was in at that time by the `ii_to_mark()` call on line 270. The test on line 271 checks for an end-of-line anchor, in which case the newline (which is part of the lexeme at this point) must be pushed back into the input (in case it is needed to match a start-of-line anchor in the next lexeme). The lexeme is terminated, and the global variables that communicate with the parser are initialized on lines 277 to 280. The `if` clause on the next

Finding the next state.

End-of-file processing,
`yywrap()`.

Failure transitions in
`yylex()`,
`yylastaccept`.

`YYBADINP`

11. I've, perhaps wrongly, perpetuated the problem in order to keep UNIX compatibility.

line removes a newline that's at the start of the lexeme as the result of a beginning-of-line anchor.

The **switch** on line 285 contains all the accepting actions that were part of the original input file. The **case** statements are all generated by LEX itself—the **case** values are the state numbers of the associated accepting state. The **default** case on line 297 should never be executed. It's here as insurance, in case an unknown bug in LEX generates a bad state number in the **switch**.

Accepting actions in `yylex()`.

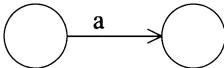
2.5 LEX—A Lexical-Analyzer Generator*

The remainder of this chapter presents the complete source code for LEX, along with the underlying theory. You must read Appendix A if you intend to look at the implementation details. The set routines presented there are used heavily in this chapter, and a familiarity with the calling conventions for these routines will be useful.

2.5.1 Thompson's Construction: From a Regular Expression to an NFA*

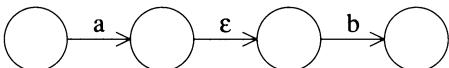
LEX constructs NFA's from regular expressions using a system called Thompson's Construction, developed by Ken Thompson at Bell Labs for the QED editor. It works as follows:

The simplest possible regular expression is a single character, and this expression can be represented by a correspondingly simple NFA. For example, a machine that matches an *a* is shown below:



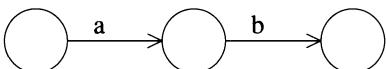
Simple expressions.

The concatenation of two regular expressions is also straightforward. The following machine represents the expression *ab* by constructing individual machines for each subexpression (the *a* and *b*), and then connecting the machines with an ϵ edge:



Concatenation.

This method needlessly wastes states, however. A better solution merges the ending state of the first machine with the start state of the second one, like this:

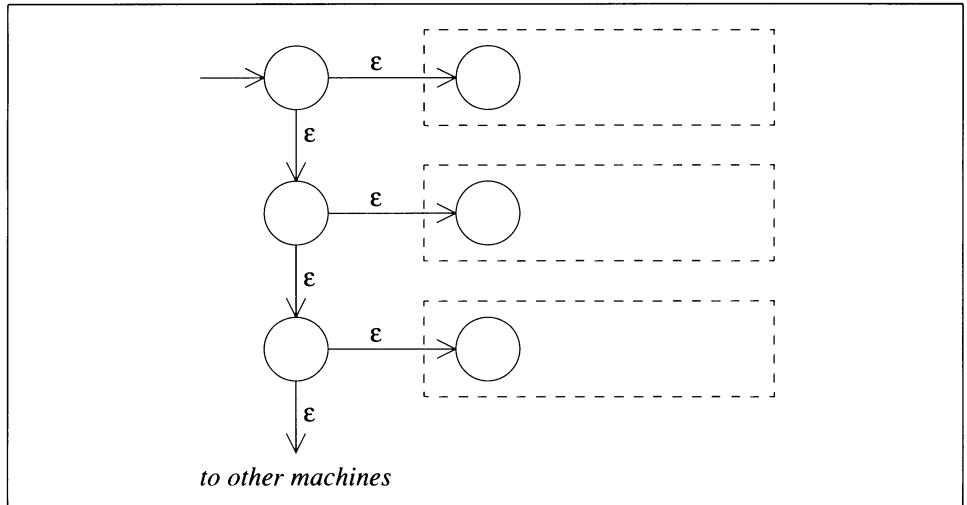


There are two situations in a LEX application where an OR of two regular expressions is required. The first is the input specification itself. That is, the LEX input contains many regular expressions, but a single machine must be output that recognizes all of these expressions. This means that all the input expressions are effectively ORed together to create the output DFA. LEX does this high-level OR using the system shown in Figure 2.9. Each of the boxes is an NFA that represents an entire regular expression, and all of these are connected together using several dummy states and ϵ edges.

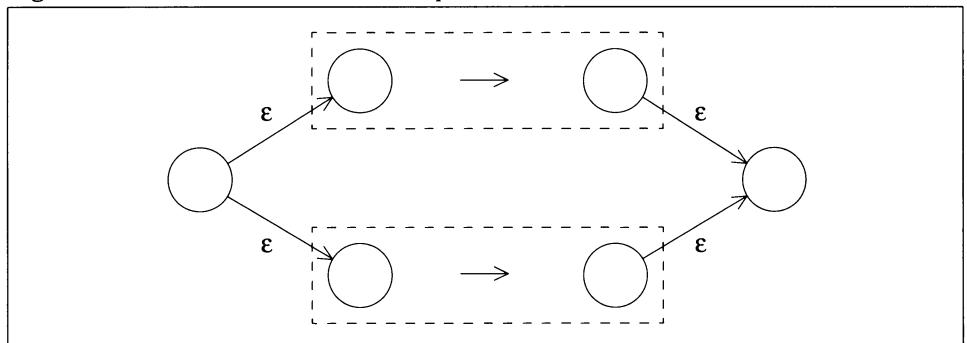
Logical OR at the top level.

The second OR situation is the OR operator (the vertical bar) which can appear in the regular expression itself (as in *alb*). LEX processes the OR operator by constructing the machine shown in Figure 2.10. Again, the seemingly empty boxes in the pictures represent machines for entire subexpressions. Figure 2.11 shows how the expression *((alb)|(cd))* would be represented. LEX starts out by making two machines to recognize the *a* and *b*, and connects the two using the OR construction shown in Figure 2.10. LEX then creates two more machines to recognize the *c* and *d*, concatenating them together by merging the end state of the first machine with the start state of the second. Finally, it

OR operator (|).

Figure 2.9. Connecting the Regular Expressions in a LEX Input File

processes the second OR operator, applying the same construction that it used earlier, but this time using the machines representing the more-complicated subexpressions (a/b) and cd in place of the boxes in Figure 2.10.

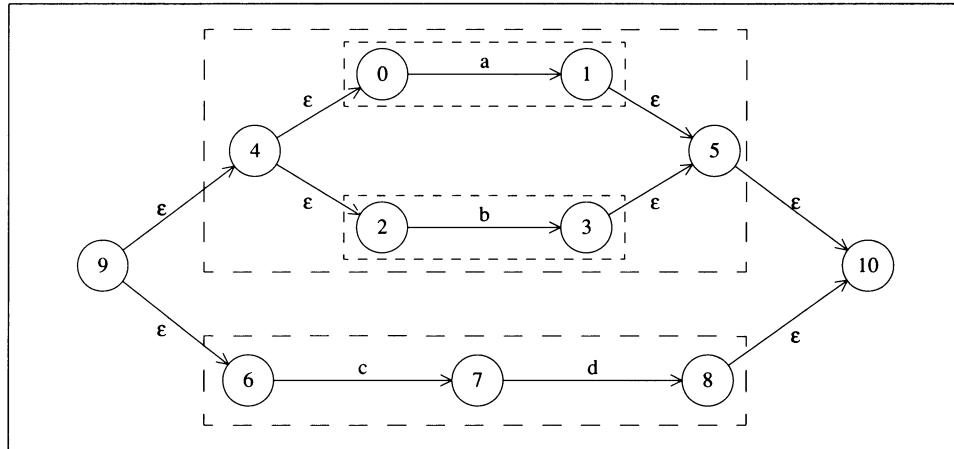
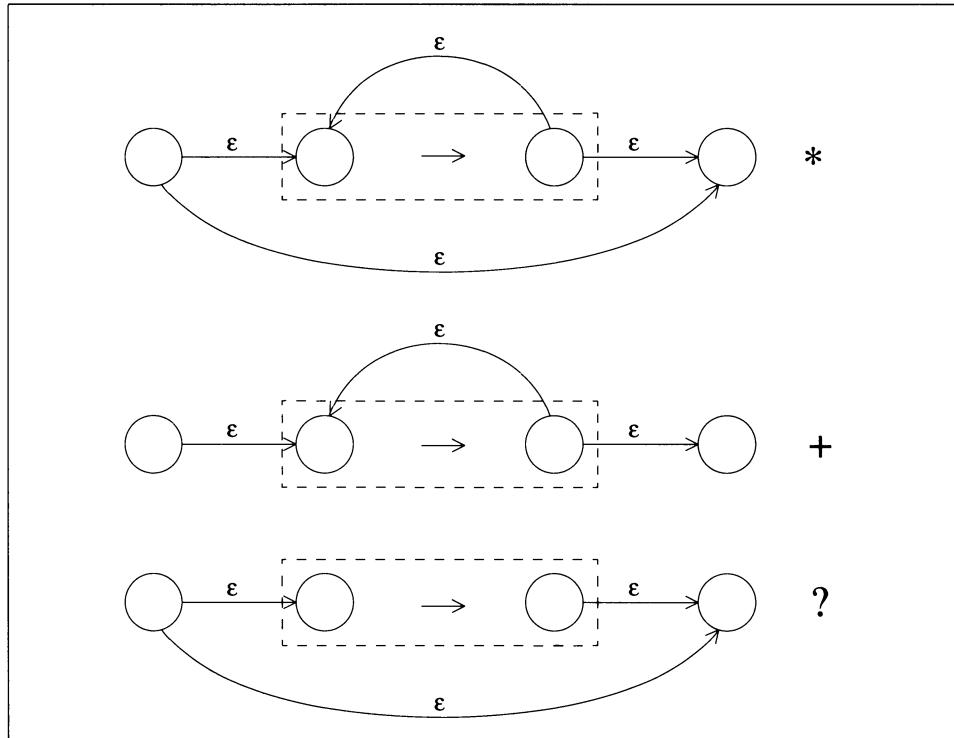
Figure 2.10. Generic NFA for the OR operator

Closure operators: * + ?

Evolution of a complex regular expression.

The machines to recognize the three closure operators are a little more complicated looking. They are shown in Figure 2.12. Note that the machines that recognize + and ? are special cases of the machine for the * operator.

Figure 2.13 shows the evolution of a machine that recognizes a subset of the floating-point constants discussed earlier. The expression used is $(D^* \cdot D | D \cdot D^*)$. It recognizes numbers with one digit to the right of the point and zero or more digits preceding it and it also recognizes the inverse—one digit to the left of the decimal point and zero or more digits following the decimal point. LEX starts out constructing an expression for the first D in the expression [in Figure 2.13(a)]. It then reads the leftmost * and substitutes the first machine into the closure machine, yielding Figure 2.13(b). It then reads the dot, and tacks it on to the right edge of the partially constructed machine [Figure 2.13(c)], and it does the same for the next D [Figure 2.13(d)]. Encountering the | operator, LEX holds onto the previously constructed expression for a moment, and then constructs a second machine for the next subexpression $(D \cdot D^*)$, not shown in the figure. Finally, it connects the two machines for the subexpressions together, using the OR construction [Figure 2.13(e)].

Figure 2.11. An NFA That Recognizes $((alb)lcd)$ **Figure 2.12.** Representing the Closure Operators

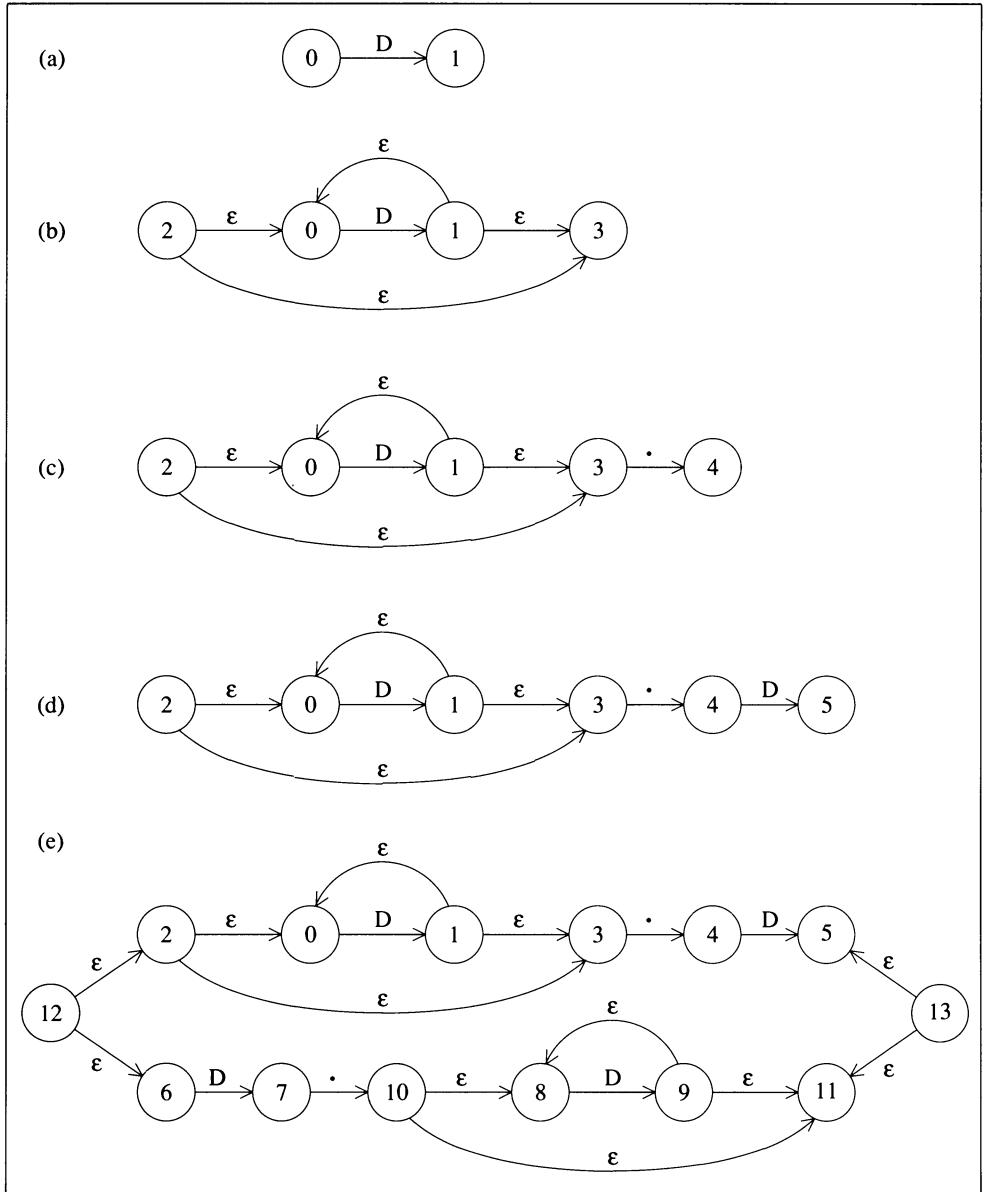
2.5.2 Implementing Thompson's Construction

This section presents a low-level implementation of the theory in the previous section. Skip forward to Section 2.5.3 if you're not interested in this level of detail.

2.5.2.1 Data Structures. A machine built with Thompson's construction has several useful characteristics:

- All the machines that recognize subexpressions—no matter how complicated—have a single start state and a single end state.
- No state has more than two outgoing edges.

Characteristics of a
Thompson machine

Figure 2.13. Constructing an NFA for $(D^*\cdot.D|D\cdot.D^*)$ 

- There are only three possibilities for the labels on the edges: (1) there is only one outgoing edge labeled with a single input character, (2) there is only one outgoing edge labeled with ϵ , and (3) there are two outgoing edges labeled with ϵ . There are never two outgoing edges labeled with input characters, and there are never two edges, one of which is labeled with an input character and the other of which is labeled with ϵ .

The NFA data structure in Listing 2.20 uses these characteristics to implement an NFA state. The next field either points at the next state, or is set to NULL if there are no outgoing edges. The next2 field is used only for states with two outgoing ϵ edges. It is set to NULL if there's only one such transition. The edge field holds one of four values that determine what the label looks like:

The NFA structure.

next, next2, and edge fields of an NFA.

- If there is a single outgoing edge labeled with an input character, edge holds that character. Single, non- ϵ edge.
- If the state has an outgoing ϵ edge, then edge is set to EPSILON. ϵ edges, EPSILON.
- If a transition is made on a character class, all characters in the class are elements of the SET pointed to by bitset, and edge holds the value CCL. The set elements are just the ASCII values of the characters. For example, if an ASCII ‘0’, which has the value 48_{10} , is in the character class, the number 48 will be in the set. Character classes, CCL, bitset.
- If the state has no outgoing transitions, edge is set to EMPTY. Terminal states, EMPTY.

Listing 2.20. *nfa.h*— Data Structures and Macros

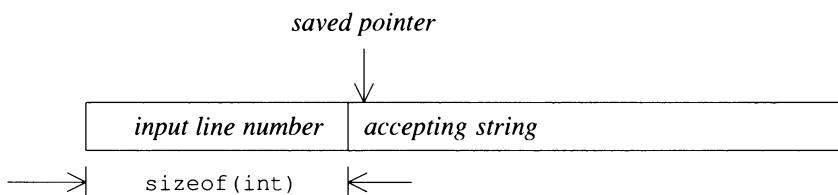
```

1  /*-----*
2   * Nfa state:
3   */
4
5  typedef struct nfa
6  {
7      int          edge;        /* Label for edge: character, CCL, EMPTY, or */
8                           /* EPSILON.                                         */
9      SET           *bitset;     /* Set to store character classes.                */
10     struct nfa  *next;       /* Next state (or NULL if none)                  */
11     struct nfa  *next2;      /* Another next state if edge==EPSILON           */
12                           /* NULL of this state isn't used               */
13     char         *accept;     /* NULL if not an accepting state, else          */
14                           /* a pointer to the action string              */
15     int          anchor;     /* Says whether pattern is anchored and, if    */
16                           /* so, where (uses #defines above).             */
17 } NFA;
18
19 #define EPSILON  -1          /* Non-character values of NFA.edge            */
20 #define CCL     -2
21 #define EMPTY    -3
22
23                           /* Values of the anchor field:                 */
24 #define NONE    0           /* Not anchored                                */
25 #define START   1           /* Anchored at start of line                  */
26 #define END     2           /* Anchored at end of line                   */
27 #define BOTH    ( START | END ) /* Anchored in both places                   */

```

The accept field in the NFA structure is NULL for nonaccepting states; otherwise, it points at a string holding the action part of the original input rule—the code to be executed when that state is accepted. The string itself is of variable length. The first **int**’s worth of characters hold the line number and the remainder of the array holds the string. For example, in a machine with a 16-bit **int**, the first two bytes of the string are the line number (in binary) and the remainder of the string is the actual input text. The pointer to the actual string is stored.

Storing accepting strings, accept.



If `p` is a pointer to an NFA structure, the actual string is accessed with `p->accept`, and the line number can be accessed by casting `accept` into a pointer to `int` and backing up one notch—like this:

```
((int*) ( p->accept ))[-1]
```

Because of various alignment problems, some care has to be taken to integrate the line number into the string in this fashion—the mechanics are discussed shortly. The technique is useful, not only in the current application, but also for doing things like attaching a precomputed string length to a string. The string can still be handled in the normal way, but the count is there when you need it.

The `anchor` field.

The macros on lines 24 to 27 of Listing 2.20 are possible values of the `anchor` field. They describe whether a `^`, `$`, or both were present in the regular expression.

`NFA_MAX`, `STR_MAX`

The remainder of `nfa.h`, in Listing 2.21, holds other definitions needed to make the machine. `NFA_MAX` is the maximum number of NFA states that can be in the machine, and `STR_MAX` is the total space available for all the accepting actions combined. The rest of the file is just prototypes for the externally-accessible functions discussed below.

Listing 2.21. `nfa.h`—Other Definitions and Prototypes

```

28 #define NFA_MAX 768           /* Maximum number of NFA states in a      */
29                      /* single machine. NFA_MAX * sizeof(NFA)   */
30                      /* can't exceed 64K.                     */
31 #define STR_MAX (10 * 1024)    /* Total space that can be used by the   */
32                      /* accept strings.                      */
33
34 void new_macro( char *definition );      /* these three are in nfa.c */
35 void printmacs( void );
36 NFA *thompson( char *(*input_funct)(), int *max_state, NFA **start_state );
37 void print_nfa( NFA *nfa, int len, NFA *start ); /* in printnfa.c */

```

Global-variable
definitions: `globals.h`

CLASS and I(x)
macros, ALLOC.

The other file you need to look at before starting is `globals.h`, in Listing 2.22. This file holds definitions for all the true global variables in the program (globals that are shared between modules). All other globals are declared `static`, so their scope is limited to the file in which they are declared. For maintenance reasons, it's desirable that both the definition and declaration of a variable be in a single file. That way you don't have to worry about maintaining two files, one in which space is allocated and the other containing `extern` statements describing the variables. The problem is solved in `globals.h` using the `CLASS` and `I(x)` macros defined on lines three to seven of Listing 2.22. The following two lines are found in only one file [typically in the same file that contains `main()`]:

```
#define ALLOC
#include "globals.h"
```

All other files include `globals.h` without the previous `ALLOC` definition. When `ALLOC` exists, `CLASS` evaluates to an empty string and `I(x)` evaluates to its argument. So, the input line:

```
CLASS char *Template I("lex.par");
```

expands to

```
char *Template = "lex.par";
```

If `ALLOC` doesn't exist, then `CLASS` expands to `extern` and `I(x)` expands to an empty string. The earlier input line expands to:

```
extern char *Template;
```

The variables on lines 11 to 15 of `globals.h` are set by command-line switches; the ones on lines 16 to 22 are used by the input routines to communicate with one another.

Listing 2.22. `globals.h`—Global-Variable Definitions

```

1  /*                               GLOBALS.H: Global variables shared between modules */
2  #ifndef ALLOC
3  #define CLASS
4  #define I(x) x
5  #else
6  #define CLASS extern
7  #define I(x)
8  #endif
9  #define MAXINP 2048           /* Maximum rule size */
10
11 CLASS int Verbose          /* Print statistics */
12 CLASS int No_lines          /* Suppress #line directives */
13 CLASS int Unix              /* Use UNIX-style newlines */
14 CLASS int Public             /* Make static symbols public */
15 CLASS char *Template        /* State-machine driver template */
16 CLASS int Actual_lineno    /* Current input line number */
17 CLASS int Lineno            /* Line number of first line of */
18                                /* a multiple-line rule. */
19 CLASS char Input_buf[MAXINP];/* Line buffer for input */
20 CLASS char *Input_file_name; /* Input file name (for #line) */
21 CLASS FILE *Ifile;          /* Input stream. */
22 CLASS FILE *Ofile;          /* Output stream. */
23 #undef CLASS
24 #undef I

```

2.5.2.2 A Regular-Expression Grammar. The code in `nfa.c`, which starts in Listing 2.23, reads a regular expression and converts it to an NFA using Thompson’s construction. The file is really a small compiler, comprising a lexical analyzer, parser, and code generator (though in this case, the generated code is a state-machine description, not assembly language). The grammar used to recognize a LEX input specification is summarized in Table 2.5. This is an informal grammar—it describes the input syntax in a general sort of way. Clarity is more important here than strict accuracy. I’ll fudge a bit in the implementation in order to get the grammar to work. Precedence and associativity are built into the grammar (the mechanics are described in depth in the next chapter). Concatenation is higher precedence than `|`; closure is higher precedence still; everything associates left to right. The various left-recursive productions have not yet been translated into an acceptable form, as was discussed in Chapter One—I’ll do that as I implement them.

2.5.2.3 File Header. The header portion of `nfa.c` is in Listing 2.23. The `ENTER` and `LEAVE` macros on lines 21 to 28 are for debugging. They expand to empty strings when `DEBUG` is not defined. When debugging, they print the current subroutine name (which is passed in as an argument), the current lexeme and what’s left of the current input line. An `ENTER` invocation is placed at the top of every subroutine of interest, and a `LEAVE` macro is put at the bottom. The text is indented by an amount proportional to the subroutine-nesting level—`Lev` is incremented by every `ENTER` invocation, and decremented by every `LEAVE`. `Lev`×4 spaces are printed to the left of every string using the `printf()`s `* field-width` capability. To simplify, the following `printf()` statement

Debugging: `ENTER`,
`LEAVE`.

Table 2.5. A Grammar for L^EX

	Productions	Notes
<i>machine</i>	$\rightarrow rule \ machine$ $\mid rule \ END_OF_INPUT$	A list of rules
<i>rule</i>	$\rightarrow expr \ EOS \ action$ $\mid ^expr \ EOS \ action$ $\mid expr \ $ \ EOS \ action$	A single regular expression followed by an accepting action. Expression anchored to start of line. Expression anchored to end of line.
<i>action</i>	$\rightarrow white_space \ string$ $\mid white_space$ $\mid \epsilon$	An optional accepting action.
<i>expr</i>	$\rightarrow expr \mid cat_expr$ $\mid cat_expr$	A list of expressions delimited by vertical bars.
<i>cat_expr</i>	$\rightarrow cat_expr \ factor$ $\mid factor$	A list of concatenated expressions.
<i>factor</i>	$\rightarrow term \ *$ $\mid term \ +$ $\mid term \ ?$ $\mid term$	A subexpression followed by a *. A subexpression followed by a +. A subexpression followed by a ?.
<i>term</i>	$\rightarrow [\ string \]$ $\mid [^ \ string \]$ $\mid [\]$ $\mid [^]$ $\mid .$ $\mid character$ $\mid (\ expr \)$	A character class. A negative character class. (nonstandard) Matches white space. (nonstandard) Matches everything but white space. Matches any character except newline. A single character. A parenthesized subexpression.
<i>white_space</i>	$\rightarrow one \ or \ more \ tabs \ or \ spaces$	
<i>character</i>	$\rightarrow any \ single \ ASCII \ character \ except \ white_space$	
<i>string</i>	$\rightarrow one \ or \ more \ ASCII \ characters$	

outputs Lev spaces by printing an empty string in a field whose width is controlled by Lev.

```
printf( "%*s", Lev, "" );
```

2.5.2.4 Error-Message Processing. The next part of *nfa.c* is the error-message routines in Listing 2.24. I've borrowed the method used by the C buffered I/O system: possible error codes are defined in the enumerated type on lines 35 to 51, and a global variable is set to one of these values when an error occurs. The *Errmsgs* array on lines 53 to 68 is indexed by error code and evaluates to an appropriate error message. Finally, the *parse_err()* subroutine on line 70 is passed an error code and prints an appropriate message. The **while** loop on line 76 tries to highlight the point at which the error occurred with a string like this:

The up arrow will (hopefully) be close to the point of error. *parse_err()* does not return.

Error messages:
Errmsgs,
parse_err().

Managing NFA structures.

new(), *discard()*.

Stack strategy,
Nfa_states[].

2.5.2.5 Memory Management. Listing 2.25 contains the memory-management routines that allocate and free the NFA structures used for the states. Two routines are used for this purpose: *new()*, on line 105, allocates a new node and *discard()*, on line 131, frees the node. I'm not using *malloc()* and *free()* because they're too slow; rather, a large array (pointed to by *Nfa_states*) is allocated the first time *new()* is called (on line 112 — the entire **if** statement is executed only once, during the first call).

A simple stack strategy is used for memory management: *discard()* pushes a pointer to the discarded node onto a stack, and *new()* uses a node from the stack if one

Listing 2.23. *nfa.c*—File Header

```

1  /* NFA.C---Make an NFA from a LeX input file using Thompson's construction */
2
3  #include <stdio.h>
4  #ifdef MSDOS
5  #   include <stdlib.h>
6  #else
7  #   include <malloc.h>
8  #endif
9
10 #include <ctype.h>
11 #include <string.h>
12 #include <tools/debug.h>
13 #include <tools/set.h>
14 #include <tools/hash.h>
15 #include <tools/compiler.h>
16 #include <tools/stack.h>
17 #include "nfa.h"                      /* defines for NFA, EPSILON, CCL */
18 #include "globals.h"                  /* externs for Verbose, etc.      */
19
20 #ifdef DEBUG
21     int    Lev = 0;
22 #   define ENTER(f) printf("%*senter %s [%c] [%1.10s] \n",
23                                Lev++ * 4, "", f, Lexeme, Input)
24 #   define LEAVE(f) printf("%*sleave %s [%c] [%1.10s]\n",
25                                --Lev * 4, "", f, Lexeme, Input)
26 #else
27 #   define ENTER(f)
28 #   define LEAVE(f)
29 #endif

```

is available, otherwise it gets a new node from the `Nfa_states[]` array (on line 124). `new()` prints an error message and terminates the program if it can't get the node. The new node is initialized with NULL pointers [the memory is actually cleared in `discard()` with the `memset()` call on line 136] and the edge field is set to `EPSILON` on line 125. The stack pointer (`Sp`) is initialized at run time on line 116 because of a bug in the Microsoft C compact model that's discussed in Appendix A.

There's an added advantage to the memory-management strategy used here. It's convenient when constructing the NFA to create a physical machine with one node per state and actual pointers to the next state. Later on, it will be convenient to have the NFA represented as an array because you can use the array index as the state number. The stack gives you both representations in a single data structure. The only disadvantage is that any nodes that are still on the stack when the NFA construction is complete will be holes in the array. It turns out that there is at most one hole, but there's no way to know in advance where it's going to be.

The other memory-management function in *nfa.c* is the string-management function, `save()`, also in Listing 2.25. This function is passed a pointer to a string, and returns a pointer to a copy—in static memory—of that string. The pointer is preceded in memory by an `int`-sized line number, as was discussed earlier on page 85. The array pointer (`Strings`, on line 100) is declared as a pointer to `int` for portability reasons. Many machines require `ints` to be aligned at more restrictive addresses than `chars`. For example, an `int` might have to be at an address that is an even multiple of four, but a `char` could be at any address. A run-time error would happen if you tried to put an `int` into an illegal address (one that was not an even multiple of four). Making `Strings` an

edge initialized to
`EPSILON`.

The same data objects
form both an array and a
graph.

String management:
`save()`.

Alignment problems
caused by leading line
number in string.

Listing 2.24. *nfa.c*—Error-Processing Routines

```

30  /*-----
31   *      Error processing stuff. Note that all errors are fatal.
32  -----
33 */
34
35 typedef enum _err_num_
36 {
37     E_MEM,           /* Out of memory */          */
38     E_BADEXPR,      /* Malformed regular expression */ */
39     E_PAREN,        /* Missing close parenthesis */ */
40     E_STACK,        /* Internal error: Discard stack full */ */
41     E_LENGTH,       /* Too many regular expressions */ */
42     E_BRACKET,      /* Missing [ in character class */ */
43     E_BOL,          /* ^ must be at start of expr or ccl */ */
44     E_CLOSE,         /* + ? or * must follow expression */ */
45     E_STRINGS,      /* Too many characters in accept actions */ */
46     E_NEWLINE,       /* Newline in quoted string */ */
47     E_BADMAC,        /* Missing } in macro expansion */ */
48     E_NOMAC,         /* Macro doesn't exist */ */
49     E_MACDEPTH,     /* Macro expansions nested too deeply. */ */
50
51 } ERR_NUM;
52
53 PRIVATE char    *Errmsgs[] =           /* Indexed by ERR_NUM */
54 {
55     "Not enough memory for NFA",
56     "Malformed regular expression",
57     "Missing close parenthesis",
58     "Internal error: Discard stack full",
59     "Too many regular expressions or expression too long",
60     "Missing [ in character class",
61     "^ must be at start of expression or after [",
62     "+ ? or * must follow an expression or subexpression",
63     "Too many characters in accept actions",
64     "Newline in quoted string, use \\n to get newline into expression",
65     "Missing } in macro expansion",
66     "Macro doesn't exist",
67     "Macro expansions nested too deeply"
68 };
69
70 PRIVATE void    parse_err( type )
71     ERR_NUM type;
72 {
73     fprintf(stderr, "ERROR (line %d) %s\n%s\n", Actual_lineno,
74                           Errmsgs[(int)type], S_input);
75
76     while( ++S_input <= Input )
77         putc('_', stderr);
78
79     fprintf( stderr, "\^\\n" );
80     exit( 1 );
81 }

```

int pointer takes care of the alignment problem at the cost of a little wasted space—the size of the region used to store the string itself must be rounded up to an even multiple of the **int** size. I’m assuming that a **char** will have less restrictive alignment rules than an **int**—a pretty safe assumption.

A single large array that will hold all the strings is allocated the first time `save()` is called on line 155 of Listing 2.25, thereby avoiding multiple inefficient `malloc()` calls every time a new string is required. The line number is put into the string on line 164, the pointer is incremented past the number, and the string itself is copied by the loop on line 166. A **char** pointer (`textp`) is used for the purpose of copying the string component. The test on line 161 is for lines starting with a vertical bar, which say that the action for the next line should be used for the current rule. No strings are copied in this situation—the same pointer is returned here as will be returned by the next `save()` call. That is, when the input string is a “|”, several pointers to the same accepting string are returned by consecutive `save()` calls. A pointer to the string itself (as compared to the line number) is returned on line 182 in order to facilitate debugging—this way you can examine the string directly without having to skip past the line number. The line number can be accessed later on, using an expression like the following:

```
char *str = save("string");
line_number = ( (int *)str )[-1];
```

Getting the line number from the left of the string.

Listing 2.25. nfa.c—Memory Management—States and String

```
82 PRIVATE NFA *Nfa_states; /* State-machine array */
83 PRIVATE int Nstates = 0; /* # of NFA states in machine */
84 PRIVATE int Next_alloc; /* Index of next element of the array */
85
86 #define SSIZE 32
87
88 PRIVATE NFA *Sstack[ SSIZE ]; /* Stack used by new() */
89 PRIVATE NFA **Sp = &Sstack[ -1 ]; /* Stack pointer */
90
91 #define STACK_OK() ( INBOUNDS(Sstack, Sp) ) /* true if stack not */
92 /* full or empty */
93 #define STACK_USED() ( (Sp-Stack) + 1 ) /* slots used */
94 #define CLEAR_STACK() ( Sp = Sstack - 1 ) /* reset the stack */
95 #define PUSH(x) ( *++Sp = (x) ) /* put x on stack */
96 #define POP() ( *Sp-- ) /* get x from stack */
97
98 /*****
99
100 PRIVATE int *Strings; /* Place to save accepting strings */
101 PRIVATE int *Savep; /* Current position in Strings array. */
102
103 ****/
104
105 PRIVATE NFA *new(); /* NFA management functions */
106 {
107     NFA *p;
108     static int first_time = 1;
109
110     if( first_time )
111     {
112         if( !( Nfa_states = (NFA *) calloc(NFA_MAX, sizeof(NFA)) ) )
113             parse_err( E_MEM );
114     }
115 }
```



Listing 2.25. continued...

```

115         first_time = 0;
116         Sp = &Sstack[ -1 ];
117     }
118
119     if( ++Nstates >= NFA_MAX )
120         parse_err( E_LENGTH );
121
122     /* If the stack is not ok, it's empty */
123
124     p = !STACK_OK() ? &Nfa_states[Next_alloc++] : POP();
125     p->edge = EPSILON;
126     return p;
127 }
128
129 /*-----*/
130
131 PRIVATE void discard( nfa_to_discard )
132 NFA *nfa_to_discard;
133 {
134     --Nstates;
135
136     memset( nfa_to_discard, 0, sizeof(NFA) );
137     nfa_to_discard->edge = EMPTY ;
138     PUSH( nfa_to_discard );
139
140     if( !STACK_OK() )
141         parse_err( E_STACK );
142 }
143
144 /*-----*/
145
146 PRIVATE char *save( str )           /* String-management function. */
147 char *str;
148 {
149     char *textp, *startp;
150     int len;
151     static int first_time = 1;
152
153     if( first_time )
154     {
155         if( !(Savep = Strings = (int *) malloc( STR_MAX )) )
156             parse_err( E_MEM );
157
158         first_time = 0;
159     }
160
161     if( *str == '|' )
162         return (char *) ( Savep + 1 );
163
164     *Savep++ = Lineno;
165
166     for( textp = (char *)Savep ; *str ; *textp++ = *str++ )
167         if( textp >= (char *)(Strings + (STR_MAX-1)) )
168             parse_err( E_STRINGS );
169
170     *textp++ = '\0' ;
171
172     /* Increment Savep past the text. "len" is initialized to the string length.
173      * The "len/sizeof(int)" truncates the size down to an even multiple of the
174      * current int size. The "+(len % sizeof(int) != 0)" adds 1 to the truncated

```



Listing 2.25. continued...

```

175     * size if the string length isn't an even multiple of the int size (the !=  

176     * operator evaluates to 1 or 0). Return a pointer to the string itself.  

177     */  

178  

179     startp = (char *)Savep;  

180     len = textp - startp;  

181     Savep += (len / sizeof(int)) + (len % sizeof(int) != 0);  

182     return startp;  

183 }

```

2.5.2.6 Macro Support. The next code segment (in Listing 2.26) comprises the macro-support routines. `new_macro()` (on line 202) is passed a pointer to a line that contains a macro definition and it files the macro in a small symbol table. The `expand_macro(char **namep)` routine on line 264 is passed a pointer to a character pointer, which in turn points at a macro invocation. The routine advances `*namep` past the invocation, and returns a string holding the macro's contents. The `printmacs()` subroutine on line 300 of Listing 2.26 prints the macros. The various hash routines and the `HASH_TAB` structure that are used here are discussed in Appendix A.

Macro substitution:
`new_macro()`,
`expand_macro()`.

`printmacs()`

Listing 2.26. nfa.c—Macro Support

```

184 -----*  

185     * MACRO support:  

186     */  

187  

188 #define MAC_NAME_MAX 34           /* Maximum name length */  

189 #define MAC_TEXT_MAX 80          /* Maximum amount of expansion text */  

190  

191 typedef struct  

192 {  

193     char name[ MAC_NAME_MAX ];  

194     char text[ MAC_TEXT_MAX ];  

195 } MACRO;  

196  

197  

198 PRIVATE HASH_TAB *Macros; /* Symbol table for macro definitions */  

199  

200 -----*  

201  

202 PUBLIC void new_macro( def )  

203     char *def;  

204 {  

205     /* Add a new macro to the table. If two macros have the same name, the  

206      * second one takes precedence. A definition takes the form:  

207      *           name <whitespace> text [<whitespace>]  

208      * whitespace at the end of the line is ignored.  

209      */  

210  

211     unsigned hash_add();  

212  

213     char *name;                  /* Name component of macro definition */  

214     char *text;                  /* text part of macro definition */  

215     char *edef;                 /* pointer to end of text part */  

216     MACRO *p;  

217     static int first_time = 1;  

218

```



Listing 2.26. continued...

```

219     if( first_time )
220     {
221         first_time = 0;
222         Macros = maketab( 31, hash_add, strcmp );
223     }
224
225     for( name = def; *def && !isspace(*def) ; def++ )           /* Isolate name */
226     ;
227     if( *def )
228         *def++ = '\0' ;
229
230     /* Isolate the definition text. This process is complicated because you need
231      * to discard any trailing whitespace on the line. The first while loop
232      * skips the preceding whitespace. The for loop is looking for end of
233      * string. If you find a white character (and the \n at the end of string
234      * is white), remember the position as a potential end of string.
235     */
236
237     while( isspace( *def ) )           /* skip up to macro body */           */
238         ++def;
239
240     text = def;                      /* Remember start of replacement text */
241
242     edef = NULL;                    /* strip trailing white space */          */
243     while( *def )
244     {
245         if( !isspace(*def) )
246             ++def;
247         else
248             for(edef = def++; isspace(*def) ; ++def )
249                 ;
250     }
251
252     if( edef )
253         *edef = '\0';
254     /* Add the macro to the symbol table */           */
255
256     p = (MACRO *) newsym( sizeof(MACRO) );
257     strncpy( p->name, name, MAC_NAME_MAX );
258     strncpy( p->text, text, MAC_TEXT_MAX );
259     addsym( Macros, p );
260 }
261
262 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
263
264 PRIVATE char    *expand_macro( namep )
265 char        **namep;
266 {
267     /* Return a pointer to the contents of a macro having the indicated
268      * name. Abort with a message if no macro exists. The macro name includes
269      * the brackets, which are destroyed by the expansion process. *namep
270      * is modified to point past the close brace.
271     */
272
273     char    *p;
274     MACRO  *mac;
275

```



Listing 2.26. continued...

```

276     if( !(p = strchr( ++(*namep), '}')) )           /* skip { and find } */
277         parse_err( E_BADMAC );
278     else
279     {
280         *p++    = '\0';                                /* Overwrite close brace. */
281
282         if( !(mac = (MACRO *) findsym( Macros, *namep )) )
283             parse_err( E_NOMAC );
284
285         *namep = p;                                     /* Update name pointer. */
286         return mac->text;
287     }
288
289     return "ERROR";                                 /* If you get here, it's a bug */
290 }
291
292 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
293
294 PRIVATE print_a_macro( MACRO *mac )                /* Workhorse function needed by      */
295 MACRO   *mac;                                      /* ptab() call in printmacs(), below */
296 {
297     printf( "%-16s--[%s]--\n", mac->name, mac->text );
298 }
299
300 PUBLIC void printmacs()                            /* Print all the macros to stdout */
301 {
302     if( !Macros )
303         printf("\tThere are no macros\n");
304     else
305     {
306         printf("\nMACROS:\n");
307         ptab( Macros, print_a_macro, NULL, 1 );
308     }
309 }
```

2.5.2.7 LEX’s Lexical Analyzer. The lowest-level input functions are in *input.c*, Listing 2.27. `get_expr()` on line eight is the actual input function. It gets an entire rule—both the regular expression and any following code—from the input file (pointed to by `Ifile`) and puts it into `Input_buf[]`. Multiple-line rules are handled here in that lines that start with white space are concatenated to the previous line. Two line-number variables are modified in this routine. They are `Lineno`, which holds the input line number of the first line of the rule, and `Actual_lineno`, which holds the current input line number. `get_expr()` normally returns a pointer to the input string (in `Input_buf[]`). It returns `NULL` either at end of file or when a line starting with a `%%` is encountered. Since `%%` is treated as an end of file, the third part of the input file, which contains C source code that is passed directly to the output, is ignored by the parser.

Listing 2.28 holds the lexical analyzer itself. The token set is defined in the enumerated type on lines 310 to 330. The `L` token (*L* for literal) is used for all characters that aren’t represented by explicitly defined tokens. Escaped characters and characters within quoted strings are also returned as `L` tokens, even if the lexeme would normally be an explicit token. The `EOS` token is returned at end of the regular expression, but the input buffer holds the entire rule, including a multiple-line accepting action. The parser uses this fact to pick up the accepting action when an `EOS` is encountered. Note that end of input is also treated as a token.

`get_expr()`

Input buffers: `Ifile`,
`Input_buffer`.

Multiple-line actions:
`Lineno`,
`Actual_lineno`.

LEX’s lexical analyzer.

Literal characters: the `L`
token, `EOS`.

Listing 2.27. *input.c*—Low-Level Input Functions

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <tools/debug.h>
4 #include "globals.h"
5
6 /* INPUT.C      Lowest-level input functions. */
7
8 PUBLIC char *get_expr()
9 {
10     /* Input routine for nfa(). Gets a regular expression and the associated
11      * string from the input stream. Returns a pointer to the input string
12      * normally. Returns NULL on end of file or if a line beginning with % is
13      * encountered. All blank lines are discarded and all lines that start with
14      * whitespace are concatenated to the previous line. The global variable
15      * Lineno is set to the line number of the top line of a multiple-line
16      * block. Actual_lineno holds the real line number.
17     */
18
19     static int lookahead = 0;
20     int space_left;
21     char *p;
22
23     p = Input_buf;
24     space_left = MAXINP;
25     if (Verbose > 1)
26         printf( "b%d: ", Actual_lineno );
27
28     if (lookahead == '%')      /* next line starts with a % sign */
29         return NULL;          /* return End-of-input marker */
30
31     Lineno = Actual_lineno ;
32
33     while( (lookahead = getline(&p, space_left-1, Ifile)) != EOF )
34     {
35         if (lookahead == 0)
36             lerror(1, "Rule too long\n");
37
38         Actual_lineno++;
39
40         if( !Input_buf[0] )           /* Ignore blank lines */
41             continue;
42
43         space_left = MAXINP - (p-Input_buf);
44
45         if( !isspace(lookahead) )
46             break;
47
48         *p++ = '\n';
49     }
50
51     if (Verbose > 1)
52         printf( "%s\n", lookahead ? Input_buf : "--EOF--" );
53
54     return lookahead ? Input_buf : NULL ;
55 }
56
57 /*-----*/
58

```



Listing 2.27. continued...

```

59  PRIVATE int getline( stringp, n, stream )
60  char      **stringp;
61  FILE      *stream;
62  {
63      /* Gets a line of input. Gets at most n-1 characters. Updates *stringp
64      * to point at the '\0' at the end of the string. Return a lookahead
65      * character (the character that follows the \n in the input). The '\n'
66      * is not put into the string.
67      *
68      * Return the character following the \n normally,
69      *      EOF at end of file,
70      *      0      if the line is too long.
71  */
72
73  static int  lookahead = 0;
74  char       *str, *startstr;
75
76  startstr = str = *stringp;
77
78  if( lookahead == 0 )                                /* initialize */
79      lookahead = getc( stream );
80
81  if( n > 0 && lookahead != EOF )
82  {
83      while( --n > 0 )
84      {
85          *str      = lookahead ;
86          lookahead = getc(stream);
87
88          if( *str == '\n' || *str == EOF )
89              break;
90          ++str;
91      }
92      *str      = '\0';
93      *stringp = str ;
94  }
95  return (n <= 0) ? 0 : lookahead ;
96 }
```

Listing 2.28. *nfa.c*—*LE*X’s own Lexical Analyzer

```

310 typedef enum token
311 {
312     EOS = 1,           /* end of string      */
313     ANY,              /* .                  */
314     AT_BOL,           /* ^                  */
315     AT_EOL,           /* $                  */
316     CCL_END,          /* ]                  */
317     CCL_START,         /* [                  */
318     CLOSE_CURLY,      /* }                  */
319     CLOSE_PAREN,       /* )                  */
320     CLOSURE,           /* *                  */
321     DASH,              /* -                  */
322     END_OF_INPUT,      /* EOF                */
323     L,                 /* literal character */
324     OPEN_CURLY,        /* {                  */
325     OPEN_PAREN,         /* (                  */
326 }
```



Listing 2.28. continued...

```

326     OPTIONAL,          /* ? */ 
327     OR,               /* | */ 
328     PLUS_CLOSE,       /* + */ 
329
330 } TOKEN;
331
332 PRIVATE TOKEN Tokmap[] =
333 {
334 /* ^@ ^A ^B ^C ^D ^E ^F ^G ^H ^I ^J ^K ^L ^M ^N */
335 L, L,
336
337 /* ^O ^P ^Q ^R ^S ^T ^U ^V ^W ^X ^Y ^Z ^[ ^\ ^] */
338 L, L,
339
340 /* ^^ ^ SPACE ! " # $ % & ' */
341 L, L, L, L, L, L, AT_EOL, L, L, L,
342
343 /* ( ) * + , - . */
344 OPEN_PAREN, CLOSE_PAREN, CLOSURE, PLUS_CLOSE, L, DASH, ANY,
345
346 /* / 0 1 2 3 4 5 6 7 8 9 : ; < = */
347 L, L,
348
349 /* > ? */
350 L, OPTIONAL,
351
352 /* @ A B C D E F G H I J K L M N */
353 L, L,
354
355 /* O P Q R S T U V W X Y Z */
356 L, L,
357
358 /* [ \ ] ^ */
359 CCL_START, L, CCL_END, AT_BOL,
360
361 /* _ ' a b c d e f g h i j k l m */
362 L, L,
363
364 /* n o p q r s t u v w x y z */
365 L, L,
366
367 /* { | } DEL */
368 OPEN_CURLY, OR, CLOSE_CURLY, L
369 };
370
371 PRIVATE char *(*Ifunct)(); /* Input function pointer */
372 PRIVATE char *Input = ""; /* Current position in input string */
373 PRIVATE char *S_input; /* Beginning of input string */
374 PRIVATE TOKEN Current_tok; /* Current token */
375 PRIVATE int Lexeme; /* Value associated with LITERAL */
376
377 #define MATCH(t) (Current_tok == (t))
378
379 /* -----
380 * Lexical analyzer:
381 *
382 * Lexical analysis is trivial because all lexemes are single-character values.
383 * The only complications are escape sequences and quoted strings, both
384 * of which are handled by advance(), below. This routine advances past the
385 * current token, putting the new token into Current_tok and the equivalent

```

Listing 2.28. continued...

```

386     * lexeme into Lexeme. If the character was escaped, Lexeme holds the actual
387     * value. For example, if a "\s" is encountered, Lexeme will hold a space
388     * character. The MATCH(x) macro returns true if x matches the current token.
389     * Advance both modifies Current_tok to the current token and returns it.
390     */
391
392 PRIVATE int      advance()
393 {
394     static int    inquote = 0;           /* Processing quoted string */
395     int          saw_esc;             /* Saw a backslash */
396     static char   *stack[ SSIZE ],    /* Input-source stack */
397                 **sp = NULL;            /* and stack pointer. */
398
399     if( !sp )                         /* Initialize sp. */
400         sp = stack - 1;              /* Necessary for large model */
401
402     if( Current_tok == EOS )        /* Get another line */
403     {
404         if( inquote )
405             parse_err( E_NEWLINE );
406         do
407         {
408             if( !(Input = (*Ifunct)()) )           /* End of file */
409             {
410                 Current_tok = END_OF_INPUT;
411                 goto exit;
412             }
413             while( isspace( *Input ) )           /* Ignore leading
414                 Input++;                      /* white space... */
415
416             } while( !*Input );                /* and blank lines. */
417
418             S_input = Input;                /* Remember start of line */
419         }                                /* for error messages. */
420
421     while( *Input == '\0' )
422     {
423         if( INBOUNDS(stack, sp) )           /* Restore previous input source */
424         {
425             Input = *sp-- ;
426             continue;
427         }
428
429         Current_tok = EOS;               /* No more input sources to restore */
430         Lexeme = '\0';                  /* ie. you're at the real end of */
431         goto exit;                     /* string. */
432     }
433
434     if( !inquote )
435     {
436         while( *Input == '{' )           /* Macro expansion required */
437         {
438             *++sp = Input ;             /* Stack current input string */
439             Input = expand_macro( sp ); /* and replace it with the macro */
440                                         /* body. */
441             if( TOOHIGH(stack,sp) )
442                 parse_err(E_MACDEPTH); /* Stack overflow */
443         }
444     }
445 }
```



Listing 2.28. continued...

```

446     if( *Input == '"' )
447     {
448         inquote = ~inquote;           /* At either start and end of a quoted      */
449         if( !*++Input )           /* string. All characters are treated as      */
450             {                      /* literals while inquote is true).          */
451                 Current_tok = EOS ;
452                 Lexeme = '\0';
453                 goto exit;
454             }
455         }
456
457     saw_esc = (*Input == '\\');
458
459     if( !inquote )
460     {
461         if( isspace(*Input) )
462         {
463             Current_tok = EOS ;
464             Lexeme = '\0';
465             goto exit;
466         }
467         Lexeme = esc( &Input );
468     }
469     else
470     {
471         if( saw_esc && Input[1] == '"' )
472         {
473             Input += 2;
474             Lexeme = '"';
475         }
476         else
477             Lexeme = *Input++ ;
478     }
479
480     Current_tok = (inquote || saw_esc) ? L : Tokmap[Lexeme] ;
481     exit:
482     return Current_tok;
483 }
```

The translation problem is simplified by the fact that all lexemes are single characters. ([^, which starts a negative character class, is treated as a **CCL_START/AT_BOL** pair). The **Tokmap[]** array on lines 332 to 369 is used to translate these single-character lexemes into tokens. It is indexed by ASCII character, and evaluates to the token value associated with that character.

Various other variables are needed by the lexical analyzer, all declared on lines 371 to 375 of Listing 2.28. **Ifunct** points at the current input function, which should work like **gets()**, returning either the next input line or **NULL** at end of file. **Input** points at the current input line, and **S_input** is the current position on that line. **Current_tok** holds the current token, and **Lexeme** holds the associated lexeme (which is only needed for **L** tokens). The **MATCH(t)** macro on line 377 of Listing 2.29 evaluates true if **t** is the current lookahead token.

The **advance()** function on lines 392 to 483 of Listing 2.28 advances the input by one character, modifying **Current_tok** and **Lexeme** as appropriate, and also returning the current token. An **END_OF_INPUT** token is returned at end of file. This subroutine is probably the most complicated routine in the parser, and is also a little long for my taste, but it seemed best to keep the whole lexical analysis phase in a single subroutine.

for maintenance reasons. A new line is fetched on lines 402 to 419 only if the end of the previous line has been reached (Current_tk is EOS).

Lines 421 to 444 of Listing 2.28 handle macro expansion. Macros are delimited by braces, and they are recognized in the `while` loop on line 436, which finds the leading brace. It’s a loop because macro definitions might be nested—if the first character of the macro body is also an open brace, the loop will expand this inner macro as well as the current one. Nested macros are handled with a stack. When an inner macro is encountered, the current input buffer is pushed onto a stack, and Input is modified to point at the macro-replacement text. The loop on line 421 is activated when you get to the end of the replacement text. The previous input string is restored at the top of the loop by popping it off the stack. The code on lines 429 to 431 is activated only if the stack is empty and no more characters are found in the current input source, in which case end of string has been reached. The `goto` statement on line 431 is functioning as a `return` statement here. I generally shun multiple `return` statements in favor of multiple `goto` branches to a label that precedes a single `return` statement. This way, the subroutine has only a single exit point so it’s easier to set up breakpoints and debugging diagnostics.

Macro expansion.

Quotes are recognized on line 446, and `inquote` is set to true when processing a quoted string. Similarly, `saw_escape` is set to true on line 457 when a backslash is encountered. The clause on lines 461 to 467 handles normal text. EOS, which marks the end of the regular expression, is returned if any white space is encountered, and escape sequences are expanded by the `esc()` call on line 467. The following `else` clause handles quoted strings. The test on line 471 is looking for a ", which must be treated as a literal quote mark. All other characters, including white space, are considered to be part of the regular expression and are just returned in the normal way. Finally, the current token is put into `Current_tk` on line 480. If you’re in a quoted string or if the current character is preceded by a backslash, then the current character is treated literally and an L token is returned; otherwise, the character is translated to a token by looking it up in `Tokmap[]`.

Quotes in regular expressions.

2.5.2.8 Parsing. LEX’s parser begins in Listing 2.29. The prototypes at the top of the listing are necessary, both for debugging and because the parser itself is highly recursive—there is no way to arrange the subroutines to avoid all forward references. The parser is a straightforward implementation of the grammar presented earlier. The topmost routine, `machine()` on line 508 of Listing 2.29, collects a series of rules and chains them together using ε transitions and dummy states (as was pictured in Figure 2.9 on page 82). The `rule()` calls on lines 516 and 522 return pointers to NFA’s that represent each regular expression on the input line.

`machine()`

The `rule()` subroutine on line 531 of Listing 2.29 gets a single regular expression and associated action. Most of the work is done by `expr()`, called on lines 554 and 557. The routine is passed two pointers to NFA pointers. When `expr()` returns, these two pointers will have been modified to point at the first and last nodes in the machine (there will be only one of each). That is, synthesized attributes are used here (and throughout the rest of the parser), but I can’t use the actual return value (the argument to a `return` statement) because there are two attributes. Consequently, I’m passing pointers to variables to be modified. (Put another way, I’m doing a call by reference here—passing pointers to the object to be modified).

`rule()`

`expr()`

Beginning- and end-of-line anchors are processed directly in `rule()` on lines 550 to 554 of Listing 2.29. An extra node is created with an outgoing transition on a newline so that the ^ is treated as if it were a regular expression that matched a newline. The anchor field is modified on line 552 to remember that this newline is there as the result of an anchor, as compared to a specific match of a \n. The newline has to be discarded in the former case; it remains in the lexeme if a specific match was requested, however.

Anchor processing.

Listing 2.29. *nfa.c*—Parser, Part 1: `machine()` and `rule()`

```

484 PRIVATE int    advance      ( void           );
485 PRIVATE void   cat_expr    ( NFA**, NFA**  );
486 PRIVATE void   discard     ( NFA*          );
487 PRIVATE void   dodash     ( SET*          );
488 PRIVATE void   expr        ( NFA**, NFA**  );
489 PRIVATE void   factor     ( NFA**, NFA**  );
490 PRIVATE int    first_in_cat ( TOKEN         );
491 PRIVATE NFA   *machine   ( void           );
492 PRIVATE NFA   *new       ( void           );
493 PRIVATE void   parse_err  ( ERR_NUM      );
494 PRIVATE NFA   *rule      ( void           );
495 PRIVATE char   *save      ( char*          );
496 PRIVATE void   term       ( NFA**, NFA**  );
497
498 /* -----
499  * The Parser:
500  *      A simple recursive descent parser that creates a Thompson NFA for
501  *      a regular expression. The access routine [thompson()] is at the
502  *      bottom. The NFA is created as a directed graph, with each node
503  *      containing pointer's to the next node. Since the structures are
504  *      allocated from an array, the machine can also be considered
505  *      as an array where the state number is the array index.
506 */
507
508 PRIVATE NFA   *machine()
509 {
510     NFA *start;
511     NFA *p;
512
513     ENTER("machine");
514
515     p = start = new();
516     p->next = rule();
517
518     while( !MATCH(END_OF_INPUT) )
519     {
520         p->next2 = new();
521         p       = p->next2;
522         p->next = rule();
523     }
524
525     LEAVE("machine");
526     return start;
527 }
528
529 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
530
531 PRIVATE NFA   *rule()
532 {
533     /* rule    --> expr  EOS action
534      *           ^expr EOS action
535      *           expr$ EOS action
536      *
537      * action  --> <tabs> <string of characters>
538      *           epsilon
539      */
540 }
```



Listing 2.29. continued...

```

541     NFA *p;
542     NFA *start = NULL;
543     NFA *end   = NULL;
544     int anchor = NONE;
545
546     ENTER("rule");
547
548     if( MATCH( AT_BOL ) )
549     {
550         start      = new() ;
551         start->edge = '\n' ;
552         anchor    |= START ;
553         advance();
554         expr( &start->next, &end );
555     }
556     else
557         expr( &start, &end );
558
559
560     if( MATCH( AT_EOL ) )
561     {
562         /* pattern followed by a carriage-return or linefeed (use a
563          * character class).
564         */
565
566         advance();
567         end->next = new()      ;
568         end->edge = CCL        ;
569
570         if( !( end->bitset = newset() ) )
571             parse_err( E_MEM );
572
573         ADD( end->bitset, '\n' );
574
575         if( !Unix )
576             ADD( end->bitset, '\r' );
577
578         end      = end->next ;
579         anchor |= END       ;
580     }
581
582     while( isspace(*Input) )
583         Input++ ;
584
585     end->accept = save( Input );
586     end->anchor = anchor;
587     advance();                      /* skip past EOS */
588
589     LEAVE("rule");
590     return start;
591 }
```

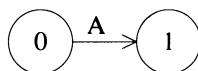
This information is eventually output as a table that will be used by the L^EX-generated driver when it processes the newline at run time. The end-of-line anchor is handled in a similar way on lines 566 to 579 of Listing 2.29, though the extra node is put at the end of the machine rather than at the beginning.

MS-DOS, end-of-line problems in anchors.

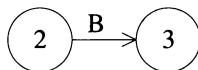
OR and concatenation:
expr(), cat_expr().

Anchors are recognized with a character class that matches either a carriage return or linefeed (as compared to a literal match of a '\n' character). You must recognize both characters in MS-DOS "binary-mode" input, because all input lines are terminated with a CR-LF pair (in that order). Lines are terminated by a single newline only in "translated" mode. Since I don't want to worry about which of the two input modes are used, I'm testing for both possibilities. When expr() returns, the input is positioned at the start of the accepting action, which is saved on line 585 (remember, an entire multiple-line action is collected into the input string).

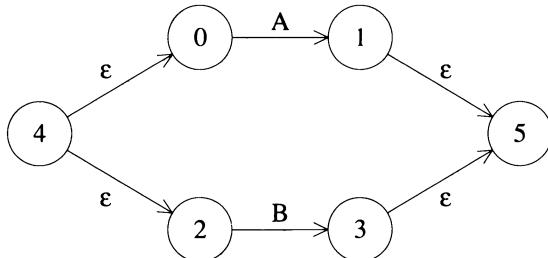
Subroutines expr() and cat_expr() are in Listing 2.30. These routines handle the binary operations: | (OR) and concatenation. I'll show how expr() works by watching it process the expression A|B. The cat_expr() call on line 621 creates a machine that recognizes the A:



*startp and *endp are modified to point at Nodes 0 and 1, and the input is advanced to the OR operator. The MATCH on line 623 succeeds and the OR is skipped by the subsequent advance() call. The second cat_expr() call (on line 625) creates a machine that recognizes the B:



and modifies e2_start and e2_end to point at nodes representing States 2 and 3. The two machines are then joined together to create the following machine:



Node 4 is created on lines 628 to 630 and *startp is modified to point at Node 4 on line 631. Node 5 is set up in a similar way on the next four lines.

Listing 2.30. nfa.c—Parser, Part 2: Binary Operators

```

592 PRIVATE void expr( startp, endp )
593 NFA      **startp, **endp ;
594 {
595     /* Because a recursive descent compiler can't handle left recursion,
596      * the productions:
597      *
598      * expr      -> expr OR cat_expr
599      *           |  cat_expr
600      *
601      * must be translated into:
602      *
603      * expr      -> cat_expr expr'
604      * expr'    -> OR cat_expr expr'
605      *           epsilon
606      *
607      * which can be implemented with this loop:
608      *
609      * cat_expr
  
```



Listing 2.30. continued...

```

610      * while( match(OR) )
611      *      cat_expr
612      *      do the OR
613      */
614
615      NFA *e2_start = NULL; /* expression to right of | */
616      NFA *e2_end   = NULL;
617      NFA *p;
618
619      ENTER("expr");
620
621      cat_expr( startp, endp );
622
623      while( MATCH( OR ) )
624      {
625          advance();
626          cat_expr( &e2_start, &e2_end );
627
628          p = new();
629          p->next2 = e2_start ;
630          p->next  = *startp ;
631          *startp = p;
632
633          p = new();
634          (*endp)->next = p;
635          e2_end ->next = p;
636          *endp = p;
637      }
638      LEAVE("expr");
639  }
640
641 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
642
643 PRIVATE void cat_expr( startp, endp )
644 NFA    **startp, **endp ;
645 {
646     /* The same translations that were needed in the expr rules are needed again
647     * here:
648     *
649     * cat_expr -> cat_expr | factor
650     *             factor
651     *
652     * is translated to:
653     *
654     * cat_expr -> factor cat_expr'
655     * cat_expr' -> | factor cat_expr'
656     *                 epsilon
657     */
658
659     NFA *e2_start, *e2_end;
660
661     ENTER("cat_expr");
662
663     if( first_in_cat( Current_tok ) )
664         factor( startp, endp );
665
666     while( first_in_cat( Current_tok ) )
667     {
668         factor( &e2_start, &e2_end );
669

```



Listing 2.30. continued...

```

670         memcpy( *endp, e2_start, sizeof(NFA) );
671         discard( e2_start );
672
673         *endp = e2_end;
674     }
675
676     LEAVE( "cat_expr" );
677 }
678
679 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
680
681 PRIVATE int      first_in_cat( TOKEN tok )
682 TOKEN tok;
683 {
684     switch( tok )
685     {
686         case CLOSE_PAREN:
687         case AT_EOL:
688         case OR:
689         case EOS:           return 0;
690
691         case CLOSURE:
692         case PLUS_CLOSE:
693         case OPTIONAL:    parse_err( E_CLOSE ); return 0;
694
695         case CCL_END:     parse_err( E_BRACKET ); return 0;
696         case AT_BOL:      parse_err( E_BOL ); return 0;
697     }
698
699     return 1;
700 }
```

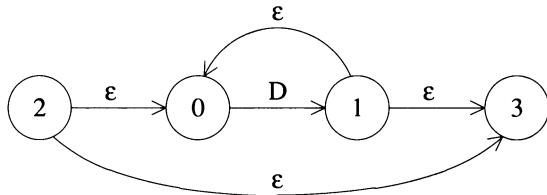
Problems with implicit concatenation operator, `first_in_cat()`.

Concatenation.

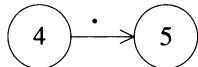
Concatenation is a somewhat harder problem because there's no operator to look for. The problem is solved by `first_in_cat()` on line 681 of Listing 2.30, which looks to see if the next input token can reasonably be concatenated to the current one. That is, there is a set of tokens that can not just be concatenated—such as the parenthesis that terminates a parenthesized subexpression or an **OR** token—and the loop must terminate when one of these is encountered. These symbols are identified by the **case** statements on lines 686 to 689. The other cases test for obvious error conditions such as a close bracket (**CCL_END**) without a preceding open bracket or one of the closure operators without anything in front of it.

Concatenation is performed in a manner similar to **OR**. The first operand is fetched with the `factor()` call on line 664 (which returns with `*startp` and `*endp` modified to point at the endpoints of a machine that recognizes the operand). The second and subsequent operands are fetched by the `factor()` call on line 668, and the two are concatenated by overwriting the contents of the end node of the first operand with the contents of the starting node of the second operand (with the `memcpy()` call on the next line). The now redundant start node of the second operand is then discarded.

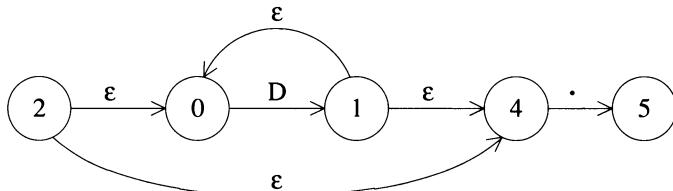
As an example of the process, if the input expression is `D* \. D`, the first `factor()` call processes the `D*`, modifying `*startp` and `*endp` to point at nodes representing States 2 and 3 of the following machine:



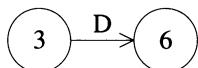
The second call modifies `e2_start` and `e2_end` to point at Nodes 4 and 5 of the following machine:



They are concatenated together by overwriting Node 3 with Node 4, yielding:



Node 3 is then discarded. The `factor()` call in the next iteration of the loop modifies `e2_start` and `e2_end` to point at the ends of:



(The discarded Node 3 is picked up again here because it will be at the top of the push-back stack.) This new machine is concatenated to the previous one by overwriting Node 5 with the contents of Node 3.

The unary closure operators are all handled by `factor()`, in Listing 2.31. It behaves just like the earlier routines, except that it builds a closure machine like the ones shown in Figure 2.12 on page 83. The code is simplified because the machines for `+` and `?` are subsets of the one for `*`. The same number of extra nodes is created in all three situations. The backwards-pointing ϵ edge is created only if a `'*'` or a `'+'` is being processed; the forward-pointing edge is created only for a `'*'` or a `'?'`.

Closure operators: `factor()`.

Single-character matches and parenthesized subexpressions are handled by `term()`, in Listing 2.32. The actual NFA structures are allocated and connected to each other on lines 761 and 762. The edge field is then initialized to the character or CCL, as appropriate. A dot is treated as a character class that matches everything but a newline. Normal character classes are assembled by `dodash()`¹² on line 811 of Listing 2.32. This routine converts the input tokens representing the class into a SET. A set can have one element (`[x]`) or several elements—`[a-zA-Z]` and the following large character class are equivalent:

`[abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ]`

Single characters and parenthesized subexpressions: `term()`.

Character classes, `dodash()`.

Note that the dash notation is expanded as sequential numbers. Since your machine probably uses the ASCII character set, this means that `[A-z]` contains the entire alphabet plus the following symbols:

12. Camptown ladies sing this song, `dodash()`, `dodash()`, Camptown race track five miles long ...

Listing 2.31. *nfa.c*—Parser, Part 3: Unary Operators (Closure)

```

701 PRIVATE void factor( startp, endp )
702 NFA    **startp, **endp;
703 {
704     /*          factor  --> term*  |  term+  |  term?
705     */
706
707     NFA *start, *end;
708
709     ENTER("factor");
710
711     term( startp, endp );
712
713     if( MATCH(CLOSURE) || MATCH(PLUS_CLOSE) || MATCH(OPTIONAL) )
714     {
715         start      = new();
716         end       = new();
717         start->next   = *startp ;
718         (*endp)->next = end      ;
719
720         if( MATCH(CLOSURE) || MATCH(OPTIONAL) )           /*  * or ? */
721             start->next2 = end;
722
723         if( MATCH(CLOSURE) || MATCH(PLUS_CLOSE) )           /*  * or + */
724             (*endp)->next2 = *startp;
725
726         *startp  = start ;
727         *endp    = end   ;
728         advance();
729     }
730
731     LEAVE("factor");
732 }
```

[\] ^ _ . ^

You can get a dash into a character class with a \-, as in [_\‐], which recognizes either a dash or an underscore. Note that the only metacharacters recognized as such in a character class are a ^ that immediately follows the bracket, a dash, and a close bracket (); so [*?.] recognizes a star, question mark, or dot. The " and { that trigger a macro expansion are also active in a character class, but they are handled by the input routines. You can use \] to get a] into a class: [[\]] recognizes either an open or close bracket.

Listing 2.32. *nfa.c*—Parser, Part 4: Single Character Matches

```

733 PRIVATE void term( startp, endp )
734 NFA    **startp, **endp;
735 {
736     /* Process the term productions:
737     *
738     * term  --> [...]  |  [^...]  |  []  |  [^]  |  .  |  (expr)  |  <character>
739     *
740     * The [] is nonstandard. It matches a space, tab, formfeed, or newline,
741     * but not a carriage return (\r). All of these are single nodes in the
742     * NFA.
743     */
744 }
```



Listing 2.32. continued...

```

745     NFA *start;
746     int c;
747
748     ENTER("term");
749
750     if( MATCH( OPEN_PAREN ) )
751     {
752         advance();
753         expr( startp, endp );
754         if( MATCH( CLOSE_PAREN ) )
755             advance();
756         else
757             parse_err( E_PAREN );
758     }
759     else
760     {
761         *startp = start           = new();
762         *endp   = start->next = new();
763
764         if( !( MATCH( ANY ) || MATCH( CCL_START) ) )
765         {
766             start->edge = Lexeme;
767             advance();
768         }
769         else
770         {
771             start->edge = CCL;
772
773             if( !( start->bitset = newset() ) )
774                 parse_err( E_MEM );
775
776             if( MATCH( ANY ) )          /*      dot (.)      */
777             {
778                 ADD( start->bitset, '\n' );
779                 if( !Unix )
780                     ADD( start->bitset, '\r' );
781
782                 COMPLEMENT( start->bitset );
783             }
784             else
785             {
786                 advance();
787                 if( MATCH( AT_BOL ) )        /* Negative character class */
788                 {
789                     advance();
790
791                     ADD( start->bitset, '\n' ); /* Don't include \n in class */
792                     if( !Unix )
793                         ADD( start->bitset, '\r' );
794
795                     COMPLEMENT( start->bitset );
796                 }
797                 if( ! MATCH( CCL_END ) )
798                     dodash( start->bitset );
799                 else                      /*      [] or [^]      */
800                     for( c = 0; c <= ' '; ++c )
801                         ADD( start->bitset, c );
802
803                     advance();
804             }
805         }
806     }

```



Listing 2.32. continued...

```

805      }
806      LEAVE ("term");
807  }
808
809 /* - - - - - - - - - - - - - - - - - - - - - - - - */
810
811 PRIVATE void      dodash( set )
812 SET      *set;                      /* Pointer to ccl character set */
813 {
814     register int      first;
815
816     for( !MATCH( EOS )   &&  !MATCH( CCL_END ) ; advance() )
817     {
818         if( ! MATCH( DASH ) )
819         {
820             first = Lexeme;
821             ADD( set, Lexeme );
822         }
823         else
824         {
825             advance();
826             for( first <= Lexeme ; first++ )
827                 ADD( set, first );
828         }
829     }
830 }
```

The final workhorse function is `printnfa()`, starting on line 57 of Listing 2.33, which is for debugging. It prints out the entire machine in human-readable form, showing the various pointers, and so forth.

Access routine: `thompson()`.

Nfa.c finishes up with a high-level access routine, `thompson()`, in Listing 2.34. (Everything else was declared PRIVATE, so was inaccessible from outside the current file). It is passed a pointer to an input function, and it returns two things: a pointer to an array of NFA structures that represents the state machine and the size of that array (the number of states in use).

Listing 2.33. *printnfa.c*— Print NFA to Standard Output

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/set.h>
4 #include <tools/hash.h>
5 #include <tools/compiler.h>
6 #include "nfa.h"
7
8 PRIVATE void      printccl      ( SET* );
9 PRIVATE char      *plab        ( NFA*, NFA* );
10 PUBLIC void      print_nfa    ( NFA*, int, NFA* );
11
12 /*****
13 * PRINTNFA.C Routine to print out a NFA structure in human-readable form.
14 */
15
16 PRIVATE void      printccl( set )
17 SET      *set;
18 {
19     static int      i;
```

Listing 2.33. continued...

```

20     putchar('[');
21     for( i = 0 ; i <= 0x7f; i++ )
22     {
23         if( TEST(set, i) )
24         {
25             if( i < ' ' )
26                 printf( "^%c", i + '@' );
27             else
28                 printf( "%c", i );
29         }
30     }
31 }
32
33     putchar(']');
34 }
35
36 /*-----*/
37
38 PRIVATE char *plab( nfa, state )
39 NFA *nfa, *state ;
40 {
41     /* Return a pointer to a buffer containing the state number. The buffer is
42      * overwritten on each call so don't put more than one plab() call in an
43      * argument to printf().
44     */
45
46     static char buf[ 32 ];
47
48     if( !nfa || !state )
49         return("--");
50
51     sprintf( buf, "%2d", state - nfa );
52     return ( buf );
53 }
54
55 /*-----*/
56
57 PUBLIC void print_nfa( nfa, len, start )
58 NFA *nfa, *start;
59 int len;
60 {
61     NFA *s = nfa ;
62
63     printf( "\n----- NFA -----\\n" );
64
65     for(; --len >= 0 ; nfa++)
66     {
67         printf( "NFA state %s: ", plab(s,nfa) );
68
69         if( !nfa->next )
70             printf("(TERMINAL)");
71         else
72         {
73             printf( "--> %s ", plab(s, nfa->next) );
74             printf( "(%s) on ", plab(s, nfa->next2) );
75
76             switch( nfa->edge )
77             {

```



Listing 2.33. continued...

```

78         case CCL:      printccl( nfa->bitset          );      break;
79         case EPSILON: printf( "EPSILON "           );      break;
80         default:     pchar   ( nfa->edge, stdout   );      break;
81     }
82 }
83
84     if( nfa == start )
85         printf(" (START STATE)");
86
87     if( nfa->accept )
88         printf(" accepting %s<%s>%s", nfa->anchor & START ? "^" : "", 
89                           nfa->accept,
90                           nfa->anchor & END   ? "$" : "" );
91         printf( "\n" );
92 }
93     printf( "\n-----\n" );
94 }
```

Listing 2.34. nfa.c— The High-Level Access Function

```

831 PUBLIC NFA      *thompson( input_function, max_state, start_state )
832 char    *(input_function)();
833 int    *max_state;
834 NFA    **start_state;
835 {
836     /* Access routine to this module. Return a pointer to a NFA transition
837      * table that represents the regular expression pointed to by expr or
838      * NULL if there's not enough memory. Modify *max_state to reflect the
839      * largest state number used. This number will probably be a larger
840      * number than the total number of states. Modify *start_state to point
841      * to the start state. This pointer is garbage if thompson() returned 0.
842      * The memory for the table is fetched from malloc(); use free() to
843      * discard it.
844 */
845
846     CLEAR_STACK();
847
848     Ifunct = input_function;
849
850     Current_tok = EOS;           /* Load first token      */
851     advance();
852
853     Nstates    = 0;
854     Next_alloc = 0;
855
856     *start_state = machine();    /* Manufacture the NFA  */
857     *max_state   = Next_alloc ; /* Max state # in NFA  */
858
859     if( Verbose > 1 )
860         print_nfa( Nfa_states, *max_state, *start_state );
861
862     if( Verbose )
863     {
864         printf("%d/%d NFA states used.\n", *max_state, NFA_MAX );
865         printf("%d/%d bytes used for accept strings.\n\n", Savep - Strings,
866                           STR_MAX );
867     }
```

Listing 2.34. continued...

```
868     return Nfa_states;
869 }
```

2.5.3 Interpreting an NFA—Theory*

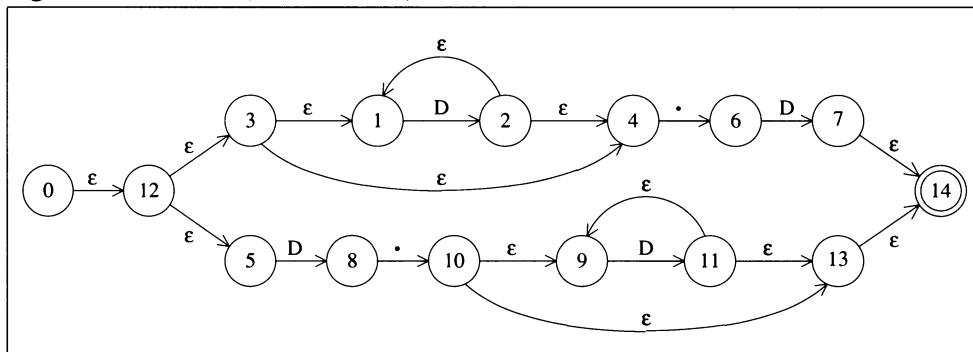
Now that we’ve constructed an NFA, we need to turn it into a DFA. The method used here is called *subset construction*; it is developed over the next few sections. First, let’s look at how an NFA can be used directly to recognize a string. (Note that interpreting an NFA directly can make more sense in some applications, such as editors, in which the time required to manufacture the DFA from the NFA can be greater than the search time.)

The basic strategy makes use of the fact that, from any state, you are, for all practical purposes, also in every state that can be made by traversing ϵ edges from the current state. I’ll demonstrate with a concrete example. The NFA for

$$(D^* \setminus D | D \setminus D^*)$$

is reproduced in Figure 2.14, and I’ll recognize the string 1 . 2. using this NFA.

Figure 2.14. NFA for $(D^* \setminus D | D \setminus D^*)$



The terminal states (the ones with no outgoing edges) are all accepting states. Starting in the start state (State 0), you can take all ϵ transitions, regardless of the input character, and without reading any input. If the machine is in State 0, it is, for all practical purposes, simultaneously in States 1, 3, 4, 5, and 12 as well. (These states are all at the end of ϵ edges that start in State 0. For example, State 4 is there by traversing three edges, from State 0 to State 12 to State 3 to State 4.) This set of states—the ones that can be reached by making ϵ transitions from a given set of start states—is called the ϵ -closure set. This set also includes the start state—there’s an implied ϵ transition from every state to itself. In this case, the set of start states has only one element: {0}, and you can make transitions to a set of five states: {1, 3, 4, 5, 12}. So:

$$\epsilon\text{-closure}(\{0\}) = \{0, 1, 3, 4, 5, 12\}.$$

Reading the 1 from the input, the next state is determined by looking at all transitions that can legitimately be made on a **D** from any of the states in the ϵ -closure set that was just created. There is a transition from State 1 to 2 on a **D**, and another from State 5 to 8; so from the states in the set {0, 1, 3, 4, 5, 12}, you can make transitions to states in the set {2, 8} on a **D**. This set of states is called the *move* set by [Aho], and it’s created by a

Using an NFA to recognize strings.

ϵ edges effectively merge states.

move() function.¹³ More formally, *move(S,c)*—where *S* is a set of states and *c* is an input character—is the set of states that can be reached by making transitions on *c* from any of the states in *S*. In this case:

$$\text{move}(\{0, 1, 3, 4, 5, 12\}, \mathbf{D}) = \{2, 8\}.$$

From States 2 and 8 the machine is also in all states that can be reached by making ϵ transitions out of States 2 and 8. There are ϵ transitions from State 2 to States 1 and 4, but there are no ϵ transitions out of State 8. So:

$$\epsilon\text{-closure}(\{2, 8\}) = \{1, 2, 4, 8\}$$

(Remember, the ϵ -closure set includes the original states.) Now get another input character, this time the dot, and look for outgoing transitions from the previously-computed ϵ -closure set. There is a transition from State 8 to 10 on a dot, and another from State 4 to 6:

$$\text{move}(\{1, 2, 4, 8\}, \cdot) = \{6, 10\}.$$

Continuing the process:

$$\epsilon\text{-closure}(\{6, 10\}) = \{6, 9, 10, 13, 14\}$$

State 14 is an accepting state, and the machine can accept if any of the states in the ϵ -closure set are accepting states. So the machine can accept at this point. Since there's more input, however, more processing is required—the same greedy algorithm that's used by the L^EX-generated analyzer is used here. Reading the 2 :

$$\text{move}(\{6, 9, 10, 13, 14\}, \mathbf{D}) = \{7, 11\}$$

and:

$$\epsilon\text{-closure}(\{7, 11\}) = \{7, 9, 11, 13, 14\}.$$

This represents an accepting state because State 14 is present. Reading the end-of-input marker,

$$\text{move}(\{7, 9, 11, 13, 14\}, \mathbf{END_OF_INPUT}) = \emptyset$$

The resulting set is empty—there's nowhere to go. This situation is the equivalent of a failure transition, so the machine executes the most recently seen accepting action and accepts State 14.

Greedy and nongreedy algorithms.

The algorithm is stated more formally in Tables 2.6 and 2.7. The first of these algorithms is a *greedy* algorithm. It recognizes the longest possible string that matches a regular expression. The second algorithm is *nongreedy*. It accepts the first string that matches a regular expression. The main difference between these algorithms is that the greedy algorithm doesn't accept an input string until a failure transition out of some state occurs, whereupon it accepts the string associated with the most recently seen accepting state (which could be the current state if the machine fails out of an accepting state). The non-greedy algorithm just accepts as soon as it enters *any* accepting state.

13. The *move()* function is the same thing as the *next()* function that we've been using to find the next state in a DFA.

Table 2.6. Greedy NFA Interpretation (Terminates on Failure)

<i>current</i>	is the set of NFA states that represents the current position in the machine.
<i>c</i>	is the current input character.
<i>accept</i>	if the most recently computed ϵ -closure set includes an accepting state, this is the state number, otherwise it's false (0).

```

last_accept = FALSE;
current = ε-closure( state_state );
while( c = nextchar( ) )
{
    if( ε-closure( move(current, c) ) ≠ ∅ )
    {
        if( accept )
            last_accept = accept ;
        current = next;
    }
    else if( last_accept )
    {
        ACCEPT( last_accept );
        last_accept = FALSE;
        current = ε-closure( state_state );
    }
    else
        ERROR;
}

```

Table 2.7. Non-Greedy NFA Interpretation (Terminates on First Accepting State)

<i>current</i>	is the set of NFA states that represents the current position in the machine.
<i>c</i>	is the current input character.
<i>accept</i>	if the most recently computed ϵ -closure set includes an accepting state, this is the state number, otherwise it's false (0).

```

current = ε-closure( state_state );
while( c = nextchar( ) )
{
    if( ε-closure( move(current, c) ) is not ∅ )
    {
        if( accept )
        {
            ACCEPT( accept );
            current = ε-closure( state_state );
        }
        else
            current = next;
    }
}

```

2.5.4 Interpreting an NFA—Implementation

The foregoing theory is all implemented by the code in *terp.c*, which starts in Listing 2.35. The `LARGEST_INT` macro on line 19 evaluates to the largest positive integer that can be represented in an `int`. It is portable—the actual size of an `int` is immaterial. It

Computing largest positive integer:
`LARGEST_INT`

Nfa, Nfa_states, nfa()
 works by shifting a number that contains all 1's to the right by one bit, shifting a 0 into the high bit. The cast to **unsigned** defeats sign extension on the right shift. Nfa is a pointer to the NFA that represents the regular expression, and Nfa_states is the number of states in that NFA. The nfa() subroutine on line 26 is passed a pointer to an input routine. It creates an NFA using the routines developed earlier, and returns the state number of the start state (its index in the array). The free_nfa() routine on line 43 discards the NFA created with the previous nfa() call.

Listing 2.35. terp.c—File Header

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <tools/debug.h>
4 #include <tools/set.h>
5 #include <tools/compiler.h>
6 #include "nfa.h"                                /* defines for NFA, EPSILON, CCL */
7
8 /-----
9 * Prototypes for subroutines in this file
10 */
11
12 PUBLIC int      nfa      ( char *(*) () );
13 PUBLIC void     free_nfa ( void );
14 PUBLIC SET      *e_closure ( SET*, char**, int* );
15 PUBLIC SET      *move     ( SET*, int );
16
17 /-----
18
19 #define LARGEST_INT (int)((unsigned)(~0)) >> 1
20
21 PRIVATE NFA   *Nfa;                            /* Base address of NFA array      */
22 PRIVATE int    Nfa_states;                     /* Number of states in NFA       */
23
24 /-----
25
26 PUBLIC int nfa( input_routine )
27 char   *(*input_routine)();
28 {
29     /* Compile the NFA and initialize the various global variables used by
30      * move() and e_closure(). Return the state number (index) of the NFA start
31      * state. This routine must be called before either e_closure() or move()
32      * are called. The memory used for the nfa can be freed with free_nfa()
33      * (in thompson.c).
34 */
35
36     NFA *sstate;
37     Nfa = thompson(input_routine, &Nfa_states, &sstate);
38     return( sstate - Nfa );
39 }
40
41 /-----
42
43 PUBLIC void free_nfa()
44 {
45     free( Nfa );
46 }
```

The ϵ -closure set is computed by `e_closure()`, in Listing 2.36. It uses the NFA created by a previous `nfa()` call. It is passed a pointer to a set of states (`input`) and returns the ϵ -closure set, which is empty if either the input set or the ϵ -closure set is empty. If the ϵ -closure set contains an accepting state, it returns two additional values indirectly through pointers—`*accept` is modified to point at the string that holds the action code, and `*anchor` is modified to hold the value of the NFA structure's anchor field. `*accept` is set to `NULL` if there is no accepting state in the closure set. Since all members of the original set are also in the closure set, the output set is created by modifying the input set. That is, `e_closure()` returns its own first argument, but the set is modified to contain the elements of the closure set in addition to the initial elements.

If the ϵ -closure set contains more than one accepting state, the accepting action that has the lowest NFA state number is used. This way, conflicting states that are higher in the input file take precedence over the ones that occur later. The `accept_num` variable declared on line 69 of Listing 2.36 holds the state number of the last-assigned accepting state. If the current state has a lower number, the other state is overwritten. It is initialized to the largest positive integer. The algorithm in Table 2.8 is used.¹⁴ The numbers in the algorithm reference comments in the code.

Computing ϵ closure:
`e_closure()`

Resolving ambiguous accepting actions. Higher actions, higher precedence

Table 2.8. Algorithm for Computing ϵ -Closure

<i>input</i>	set of input states.
<i>i</i>	the current NFA state being examined.
<i>N</i>	state number of a next state that can be reached from State <i>i</i> .
1:	Push all states in the input set onto a stack.
2:	while(the stack is not empty)
	{
3:	Pop the top element into <i>i</i> .
	if(State <i>i</i> is an accepting state)
	*accept = the accept string;
4:	if(there's an ϵ transition from State <i>i</i> to State <i>N</i>)
	{
5:	if(<i>N</i> isn't in the closure set)
	{
6:	Add <i>N</i> to the closure set.
7:	Push <i>N</i> onto the stack.
	}
	}
	}

The `move` set is figured with the `move()` subroutine in Listing 2.37. It returns either the set of states that can be reached by making transitions on a specified input character from a specified set of states, or `NULL` if there are no such transitions.

Computing the move set:
`move()`.

The `for` loop on lines 129 to 143 does the work. The `if` statement at the top of the loop evaluates true for every element of the input set, and any next states are added to the output set on line 140. Note that the output set is not created unless there's something to put into it.

14. It is derived from the one in [Aho] p. 119.

Listing 2.36. *terp.c*—The ϵ -Closure Function

```

47 PUBLIC SET *e_closure( input, accept, anchor )
48 SET    *input ;
49 char   **accept ;
50 int    *anchor ;
51 {
52     /* input      is the set of start states to examine.
53      * *accept  is modified to point at the string associated with an accepting
54      *           state (or to NULL if the state isn't an accepting state).
55      * *anchor   is modified to hold the anchor point, if any.
56      *
57      * Computes the epsilon closure set for the input states. The output set
58      * will contain all states that can be reached by making epsilon transitions
59      * from all NFA states in the input set. Returns an empty set if the input
60      * set or the closure set is empty, modifies *accept to point at the
61      * accepting string if one of the elements of the output state is an
62      * accepting state.
63     */
64
65     int stack[ NFA_MAX ];          /* Stack of untested states      */
66     int *tos;                     /* Stack pointer                  */
67     NFA *p;                      /* NFA state being examined     */
68     int i;                        /* State number of "            */
69     int accept_num = LARGEST_INT ;
70
71     if( !input )
72         goto abort;
73
74     *accept = NULL;               /* Reference to algorithm: */   /* 1 */
75     tos     = & stack[-1];
76
77     for( next_member(NULL); (i = next_member(input)) >= 0 ;)
78         *++tos = i;
79
80     while( INBOUNDS(stack,tos) )           /* 2 */
81     {
82         i = *tos-- ;
83         p = & Nfa[ i ];
84         if( p->accept && (i < accept_num) )
85         {
86             accept_num = i ;
87             *accept    = p->accept ;
88             *anchor    = p->anchor ;
89         }
90
91         if( p->edge == EPSILON )           /* 4 */
92         {
93             if( p->next )
94             {
95                 i = p->next - Nfa;
96                 if( !MEMBER(input, i) )           /* 5 */
97                 {
98                     ADD( input, i );           /* 6 */
99                     *++tos = i;              /* 7 */
100                }
101            }
102            if( p->next2 )
103            {
104                i = p->next2 - Nfa;

```



Listing 2.36. continued...

```

105             if( !MEMBER(input, i) )                                /* 5 */
106             {
107                 ADD( input, i );
108                 *++tos = i;
109             }
110         }
111     }
112 }
113 abort:
114     return input;
115 }
```

Listing 2.37. *terp.c*—The Move Function

```

116 PUBLIC SET *move( inp_set, c )
117 SET    *inp_set;                                     /* input set           */
118 int    c;                                         /* transition on this character */
119 {
120     /* Return a set that contains all NFA states that can be reached by making
121      * transitions on "c" from any NFA state in "inp_set." Returns NULL if
122      * there are no such transitions. The inp_set is not modified.
123      */
124
125     int i;
126     NFA *p;                                         /* current NFA state   */
127     SET *outset = NULL;                            /* output set          */
128
129     for( i = Nfa_states; --i >= 0; )
130     {
131         if( MEMBER(inp_set, i) )
132         {
133             p = &Nfa[i];
134
135             if( p->edge==c || (p->edge==CCL && TEST(p->bitset, c)) )
136             {
137                 if( !outset )
138                     outset = newset();
139
140                 ADD( outset, p->next - Nfa );
141             }
142         }
143     }
144     return( outset );
145 }
```

Listing 2.38 uses the previous subroutines to build a small **egrep**-like program that copies standard input to standard output, printing only those lines that contain a match of a regular expression passed into the program on the command line. The `next_char()` subroutine on line 156 of Listing 2.38 gets the next input character. It buffers entire lines because `terp` must print the entire line when a match is found [otherwise `getc()` could be used]. `getline()` on line 177 is the input function passed to `nfa()`. It just passes along the regular expression that was on the command line. (`Expr` is initialized on line 215.) The remainder of the subroutine is a straightforward implementation of the non-greedy algorithm discussed in the previous section.

Listing 2.38. *terp.c*—A Simplified Egrep

```

146 #ifdef MAIN
147 #define ALLOCATE
148 #include "globals.h" /* externs for Verbose */
149
150 #define BSIZE 256
151
152 PRIVATE char Buf[ BSIZE ]; /* input buffer */
153 PRIVATE char *Pbuf = Buf; /* current position in input buffer */
154 PRIVATE char *Expr; /* regular expression from command line */
155
156 int nextchar()
157 {
158     if( !*Pbuf )
159     {
160         if( !fgets(Buf, BSIZE, stdin) )
161             return NULL;
162         Pbuf = Buf;
163     }
164     return *Pbuf++;
165 }
166
167 /-----*/
168
169 PRIVATE void printbuf()
170 {
171     fputs(Buf, stdout); /* Print the buffer and force a read */
172     *Pbuf = 0; /* on the next call to nextchar(). */
173 }
174
175 /-----*/
176
177 PRIVATE char *getline()
178 {
179     static int first_time_called = 1;
180
181     if( !first_time_called )
182         return NULL;
183
184     first_time_called = 0;
185     return Expr;
186 }
187
188 /-----*/
189
190 main( argc, argv )
191 char **argv;
192 {
193     int sstate; /* Starting NFA state */
194     SET *start_dfastate; /* Set of starting nfa states */
195     SET *current; /* current DFA state */
196     SET *next;
197     int accept; /* cur. DFA state is an accept */
198     int c; /* current input character */
199     int anchor;
200
201     if( argc == 2 )
202         fprintf(stderr,"expression is %s\n", argv[1] );

```



Listing 2.38. continued...

```

203     else
204     {
205         fprintf(stderr,"usage: terp pattern <input\n");
206         exit(1);
207     }
208
209     /* 1: Compile the NFA; initialize move() & e_closure().
210      * 2: Create the initial state, the set of all NFA states that can
211      *      be reached by making epsilon transitions from the NFA start state.
212      * 3: Initialize the current state to the start state.
213 */
214
215     Expr    = argv[1];                      /* 1 */
216     sstate = nfa( getline );
217
218     next = newset();                      /* 2 */
219     ADD( next, sstate );
220     if ( !(start_dfastate = e_closure(next, &accept, &anchor)) )
221     {
222         fprintf(stderr, "Internal error: State machine is empty\n");
223         exit(1);
224     }
225
226     current = newset();                  /* 3 */
227     assign( current, start_dfastate );
228
229     /* Now interpret the NFA: The next state is the set of all NFA states that
230      * can be reached after we've made a transition on the current input
231      * character from any of the NFA states in the current state. The current
232      * input line is printed every time an accept state is encountered.
233      * The machine is reset to the initial state when a failure transition is
234      * encountered.
235 */
236
237     while( c = nextchar() )
238     {
239         if ( next = e_closure( move(current, c), &accept, &anchor) )
240         {
241             if ( accept )
242                 printbuf();                      /* accept */
243             else
244             {                                /* keep looking */
245                 delset( current );
246                 current = next;
247                 continue;
248             }
249         }
250
251         delset( next );                  /* reset */
252         assign( current, start_dfastate );
253     }
254 }
255 #endif

```

Subset construction: an example.

2.5.5 Subset Construction: Converting an NFA to a DFA—Theory*

You can apply the foregoing procedure to translate NFA's to DFA's by computing the ϵ -closure and move sets for every possible input character. This generalized method, called *subset construction*, takes advantage of the fact that all NFA states connected with ϵ edges are effectively the same state. A single DFA state represents all NFA states that are connected in this manner. The outgoing edges from the DFA state are just the sum of the outgoing edges from all states in the ϵ -closure set. I'll demonstrate with an example. Using the NFA in Figure 2.14 on page 113, the starting DFA state consists of the ϵ -closure of the starting NFA state:

$$\epsilon\text{-closure}(\{12\}) = \{0, 1, 3, 4, 5, 12\} \quad (\text{new DFA State } 0)$$

The next states in the DFA are then computed by figuring $\text{move}(\text{current_state}, c)$, for every possible input character. The input *alphabet* (the set of all possible input characters) contains only two elements: **D** and dot (.), so DFA States 1 and 2 are computed as follows:

$$\begin{aligned} \text{move}(\{0, 1, 3, 4, 5, 12\}, \mathbf{D}) &= \{2, 8\} & (\text{new DFA State 1}) \\ \text{move}(\{0, 1, 3, 4, 5, 12\}, .) &= \{6\} & (\text{new DFA State 2}) \end{aligned}$$

This procedure is then applied to the two new states. Starting with State 2:

$$\begin{aligned} \text{DFA State 2} &= \{6\} \\ \epsilon\text{-closure}(\{6\}) &= \{6\} \\ \text{move}(\{6\}, .) &= \emptyset \\ \text{move}(\{6\}, \mathbf{D}) &= \{7\} & (\text{new DFA State 3}) \end{aligned}$$

$$\begin{aligned} \text{DFA State 1} &= \{2, 8\} \\ \epsilon\text{-closure}(\{2, 8\}) &= \{1, 2, 4, 8\} \\ \text{move}(\{1, 2, 4, 8\}, \mathbf{D}) &= \{2\} & (\text{new DFA State 4}) \\ \text{move}(\{1, 2, 4, 8\}, .) &= \{6, 10\} & (\text{new DFA State 5}) \end{aligned}$$

If the move operation results in an empty set—as is the case with $\text{move}(\{6\}, .)$, above—then there are no outgoing transitions from that state. Applying the procedure again to the three new DFA States:

$$\begin{aligned} \text{DFA State 3} &= \{7\} \\ \epsilon\text{-closure}(\{7\}) &= \{7, 13, 14\} \\ \text{move}(\{7, 13, 14\}, .) &= \emptyset \\ \text{move}(\{7, 13, 14\}, \mathbf{D}) &= \emptyset \end{aligned}$$

$$\begin{aligned} \text{DFA State 4} &= \{2\} \\ \epsilon\text{-closure}(\{2\}) &= \{1, 2, 4\} & (\text{existing DFA State 4}) \\ \text{move}(\{1, 2, 4\}, .) &= \{6\} & (\text{existing DFA State 2}) \\ \text{move}(\{1, 2, 4\}, \mathbf{D}) &= \{2\} & (\text{existing DFA State 4}) \end{aligned}$$

$$\begin{aligned} \text{DFA State 5} &= \{6, 10\} \\ \epsilon\text{-closure}(\{6, 10\}) &= \{6, 9, 10, 13, 14\} \\ \text{move}(\{6, 9, 10, 13, 14\}, .) &= \emptyset \\ \text{move}(\{6, 9, 10, 13, 14\}, \mathbf{D}) &= \{7, 11\} & (\text{new DFA State 6}) \end{aligned}$$

DFA States 3 and 5 are accepting states because their ϵ -closure sets all contain an accepting state (NFA State 14). This last expansion introduced only one new state, which is now expanded:

DFA State 6 = {7, 11}
 ϵ -closure({7, 11}) = {7, 9, 11, 13, 14}
 $\text{move}(\{7, 9, 11, 13, 14\}, .) = \emptyset$
 $\text{move}(\{7, 9, 11, 13, 14\}, D) = \{11\}$ (new DFA State 7)

And then this new state is expanded in turn:

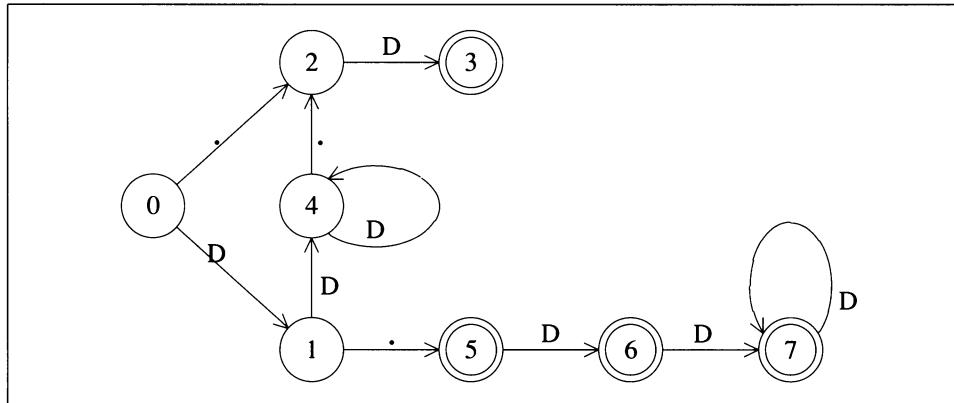
DFA State 7 = {11}
 ϵ -closure({11}) = {9, 11, 13, 14}
 $\text{move}(\{9, 11, 13, 14\}, .) = \emptyset$
 $\text{move}(\{9, 11, 13, 14\}, D) = \{11\}$ (existing DFA State 7)

DFA States 6 and 7 are accepting states because they contain NFA accepting State 14. The process is summarized in Table 2.9, and the complete DFA is pictured in Figure 2.15. The procedure is formalized with the algorithm in Table 2.10.¹⁵ Note that the state machine that this algorithm generates is not optimal. It has several more states than necessary. I'll look at how to remove these extra states in a moment.

Table 2.9. Converting NFA in Figure 2.14 (page 113) to a DFA

DFA State	NFA State	ϵ -closure	move(D)		move(.)		Accepting
			NFA	DFA	NFA	DFA	
0	{0}	{0,1,3,4,5,12}	{2,8}	1	{6}	2	no
1	{2,8}	{1,2,4,8}	{2}	4	{6, 10}	5	no
2	{6}	{6}	{7}	3	\emptyset	—	no
3	{7}	{7,13,14}	\emptyset	—	\emptyset	—	yes
4	{2}	{1,2,4}	{2}	4	{6}	2	no
5	{6, 10}	{6,9,10,13,14}	{7,11}	6	\emptyset	—	yes
6	{7, 11}	{7,9,11,13,14}	{11}	7	\emptyset	—	yes
7	{11}	{9,11,13,14}	{11}	7	\emptyset	—	yes

Figure 2.15. DFA for NFA in Figure 2.14



15. Adapted from [Aho], p. 118.

Table 2.10. Algorithm to Convert an NFA to a DFA

<i>Dstates</i>	is an array of DFA states. Each state is represented by a <u>set</u> of NFA states and a Boolean <i>mark</i> field.
<i>Dtran</i>	is the DFA transition matrix, an array as wide as the input character set and as deep as the maximum number of states. $Dtran[current_state][input_character] =$ the next state
<i>i</i>	is the index in <i>Dtran</i> of the current state.
<i>nstates</i>	is the index in <i>Dtran</i> of the place where a new state will be inserted.
<i>S</i>	is a set of NFA states that defines a DFA state.
<i>c</i>	is a possible input symbol.
<i>Initially:</i>	<p><i>Dstates[0].set</i> = ϵ-closure of NFA start state <i>Dstates[0]</i> is not marked. <i>Nstates</i> = 1. <i>i</i> = 0. <i>S</i> is empty. <i>Dtran</i> is initialized to all failure transitions.</p> <pre> while(there is an unmarked state at <i>Dstates[i]</i>) { <i>Dstates[i].mark</i> = TRUE ; for(each input symbol <i>c</i>) { <i>S</i> = ϵ-closure(move(<i>Dstate[i].set</i>, <i>c</i>)); if(<i>S</i> is not \emptyset) { if(a state with <i>set=S</i> isn't in <i>Dstates</i>) add a new unmarked state at <i>Dstates[Nstates]</i>, with <i>set=S</i> <i>Dtran[i][c]</i> = <i>Nstates++</i> else <i>Dtran[i][c]</i> = index of exiting state in <i>Dstates</i> } } } </pre>

2.5.6 Subset Construction: Converting an NFA to a DFA—Implementation

This section shows how LEX applies the theory in the previous section to make the output DFA. It uses the `e_closure()` and `move()` functions to build a complete DFA transition matrix from a previously constructed NFA. As before, you need to start with some constant and data-structure definitions, found in `dfa.h`, Listing 2.39.

The maximum number of DFA states is defined on line seven of Listing 2.40 to be 254. This number is more than adequate for most applications and has the added advantage of fitting into an **unsigned char**. The tables used for the transition matrix can be smaller as a consequence. You can make the number larger if you wish, but you'll have to change the driver discussed at the beginning of this chapter to compensate for the change in type. The **TTYPE** definition on line 12 is the type of the output tables—it must match the declared type of the tables in the state-machine driver. The internal versions of the same tables are all arrays of **int**. The **F** macro, defined on line 18, marks failure transitions in the internal tables. It must be -1, but hiding the -1 in a macro makes the code a little more readable. The **ROW** type on line 21 is one row of the DFA transition matrix. Declaring a row as a distinct type helps you use pointers to traverse individual

Output-table type: TTYPE

Failure transitions: F

The BOW type

Listing 2.39. *dfa.h*—Definitions for a DFA

```

1  /*-
2   * DFA.H: The following definitions are used in dfa.c and in
3   * minimize.c to represent DFA's.
4   *-----
5   */
6
7  #define DFA_MAX      254          /* Maximum number of DFA states. If this
8                               * number >= 255, you'll have to change the
9                               * output routines and driver. States are
10                              * numbered from 0 to DFA_MAX-1
11                              */
12 typedef unsigned char TTYPE;    /* This is the type of the output DFA      */
13                                         /* transition table (the internal one is  */
14                                         /* an array of int). It is used only to   */
15                                         /* figure the various table sizes printed */
16                                         /* by -V.                                */
17
18 #define F           -1          /* Marks failure states in the table.     */
19 #define MAX_CHARS    128         /* Maximum width of dfa transition table. */
20
21 typedef int ROW[ MAX_CHARS ];  /* One full row of Dtran, which is itself */
22                                         /* an array, DFA_MAX elements long, of    */
23                                         /* ROWs.                                 */
24 /*-----*/
25
26 typedef struct accept
27 {
28     char *string;        /* Accepting string; NULL if nonaccepting. */
29     int anchor;          /* Anchor point, if any. Values are defined in NFA.H. */
30
31 } ACCEPT;
32
33 /*-
34  * External subroutines:
35  */
36
37 SET *e_closure (SET*, char **, int *); /* terp.c      */
38 void free_nfa (void);                  /* terp.c      */
39 SET *move (SET*, int);                /* terp.c      */
40 int nfa (char *(*)());               /* terp.c      */
41 int dfa (char *(*)(), ROW*[], ACCEPT**); /* dfa.c       */
42 int min_dfa (char *(*)(), ROW*[], ACCEPT**); /* minimize.c */
43 int ttypesize (int);                 /* printdfa.c */
44 void pdefines (FILE*, int);          /* pdefines.c */
45 int columns (FILE*, ROW*, int, int, char*); /* columns.c  */
46 void cnext (FILE*, char*);          /* columns.c  */
47
48 void pheader(FILE *fp,    ROW dtran[], int nrows, ACCEPT *accept); /* print.c */
49 void pdriver(FILE *output,           int nrows, ACCEPT *accept); /* print.c */

```

columns, because incrementing a pointer to an entire row will skip as many array elements as are in the row, effectively moving you down to the next element of the current column. For example, the following code allocates a NUM_ROWS element array (a two-dimensional array of size MAX_CHARS × NUM_ROWS) and then prints the sixth column. It initializes the column pointer (*p*) to the head of the sixth column, and moves the pointer to the next element of the same column with a simple increment on each

iteration of the loop. Note that two stars are needed to get an array element in the `printf()` statement.¹⁶

```
#define NUM_ROWS 10

static ROW Array[ NUM_ROWS ];
ROW      *p;
int       i;

...
p = (ROW *) ( &Array[0][6] ); /* Initialize p to head */
                                /* of column 6.           */

for( i = NUM_ROWS; --i >= 0; printf("%d ", **p++) )
;
```

ACCEPT structures

DFA_STATE: accept,
group, anchor, set
fields

The transition matrix:
Dtran, Nstates
Create a DFA: dfa().

Returning to Listing 2.39, the ACCEPT structure on line 26 keeps track of the accepting strings, which are copied to an array of ACCEPT structures so that the memory used for the NFA can be discarded once it is translated into a DFA. The remainder of the file is just function prototypes.

The actual conversion routines start in Listing 2.40 with a few more definitions. The DFA_STATE structure, `typedefed` on lines 23 to 31, represents a single state in the DFA. The group field is used for minimization, discussed in the next section. The accept and anchor fields are the accepting string and anchor point copied from the NFA state in the case of an accepting state; set is the set of NFA states that comprise the current DFA state. Dstates, an array of DFA structures, is defined on line 33. The DFA state number is the array index. The actual transition matrix Dtran is declared on line 36, and the number of states Nstates is declared on the next line.

The routines that create the DFA are all local to *dfa.c* and are accessed via the `dfa()` routine, also in Listing 2.40 (on It is passed an input function, and returns the number of states in the transition matrix, the matrix itself, and an array holding the accepting string (the last two are returned indirectly through pointers). For the most part, this function is just allocating memory for the new tables and discarding memory for other tables once they're no longer needed. The NFA is constructed on line 69; it's converted to a DFA on line 81; the accepting strings are copied to the target array on line 91; and if verbose mode is active, a diagnostic message is printed on lines 101 to 114 of Listing 2.40.

16. C note: You need two stars because *p* points at a two-dimensional array. To see why, in C the square brackets can always be interpreted with the following identity: $p[i] \equiv * (p+i)$. If *i* is 0, then $p[0] \equiv *p$. Applying this identity to a two-dimensional array: $p[0][0] \equiv (*p)[0] \equiv *p$. Looked at another way, if *p* is a pointer to an array, then $*p$ must be the array itself. But arrays are always represented in C by a pointer to the first element. Consequently, if *p* is a pointer to an array of `int`, then $*p$ must be an array of `int`, which is represented by a pointer to the first element—a pointer to an `int`. Both *p* and $*p$ will hold the same address, but *p* is of type “pointer to array of `int`” and $*p$ is of type “pointer to `int`.”

Using a pointer here is much more efficient than something like this:

```
for( i = 0; i < NUM_ROWS; i++ )
    printf("%d ", Array[i][6] )
```

because the array index forces the compiler to multiply *i* by the size of a row at every iteration of the loop. The following is better, but still not as good as the original version:

```
ROW *p;
for( p=Array, i=NUM_ROWS; --i >= 0; printf("%d ", (*p++)[6] )
;
```

Listing 2.40. *dfa.c*—NFA to DFA Conversion: File Header and Access Function

```

1 #include <stdio.h>
2 #ifdef MSDOS
3 # include <stdlib.h>
4 #else
5 # include <malloc.h>
6 #endif
7 #include <tools/debug.h>
8 #include <tools/set.h>
9 #include <tools/compiler.h>
10 #include "dfa.h"
11 #include "globals.h"           /* externs for Verbose, etc. */
12
13 /-----
14 * DFA.C   Make a DFA transition table from an NFA created with
15 *           Thompson's construction.
16 *-----
17 * Dtran is the deterministic transition table. It is indexed by state number
18 * along the major axis and by input character along the minor axis. Dstates
19 * is a list of deterministic states represented as sets of NFA states.
20 * Nstates is the number of valid entries in Dtran.
21 */
22
23 typedef struct dfa_state
24 {
25     unsigned group : 8;          /* Group id, used by minimize() */ *
26     unsigned mark : 1;          /* Mark used by make_dtran() */ *
27     char    *accept;           /* accept action if accept state */ *
28     int    anchor;            /* Anchor point if an accept state */ *
29     SET      *set;             /* Set of NFA states represented by */ *
30
31 } DFA_STATE;                  /* this DFA state */ *
32
33 PRIVATE DFA_STATE *Dstates ; /* DFA states table */ *
34 /-----*/ *
35
36 PRIVATE ROW    *Dtran ;      /* DFA transition table */ *
37 PRIVATE int    Nstates ;     /* Number of DFA states */ *
38 PRIVATE DFA_STATE *Last_marked; /* Most-recently marked DFA state in Dtran */ *
39
40 extern char    *bin_to_ascii ( int, int ); /* in compiler.lib */
41
42 /-----
43 * Prototypes for subroutines in this file:
44 */
45
46 PRIVATE DFA_STATE *get_unmarked ( void );           ;
47 PRIVATE int      in_dstates ( SET* );               ;
48 PRIVATE void    free_sets ( void );                 ;
49 PRIVATE void    make_dtran ( int );                 ;
50 PUBLIC int     dfa ( char **(*)(), ROW*, ACCEPT** );
51 /-----*/ *
52
53 int    dfa( ifunct, dfap, acceptp )
54 char  *( *ifunct )();
55 ROW    *( dfap[] );
56 ACCEPT *( *acceptp );
57 {

```



Listing 2.40. continued...

```

58     /* Turns an NFA with the indicated start state (sstate) into a DFA and
59      * returns the number of states in the DFA transition table. *dfap is
60      * modified to point at that transition table and *acceptp is modified
61      * to point at an array of accepting states (indexed by state number).
62      * dfa() discards all the memory used for the initial NFA.
63      */
64
65     ACCEPT      *accept_states;
66     int          i;
67     int          start;
68
69     start      = nfa( ifunct );                                /* make the nfa */
70     Nstates    = 0;
71     Dstates    = (DFA_STATE *) calloc( DFA_MAX, sizeof(DFA_STATE) );
72     Dtran      = (ROW *)           calloc( DFA_MAX, sizeof(ROW) );
73     Last_marked = Dstates;
74
75     if( Verbose )
76         fputs("making DFA: ", stdout);
77
78     if( !Dstates || !Dtran )
79         ferr( "Out of memory!" );
80
81     make_dtran( start );           /* convert the NFA to a DFA           */
82     free_nfa();                  /* Free the memory used for the nfa   */
83                               /* itself (but not the accept strings). */
84
85     Dtran      = (ROW *) realloc( Dtran, Nstates * sizeof(ROW) );
86     accept_states = (ACCEPT*) malloc (      Nstates * sizeof(ACCEPT));
87
88     if( !accept_states || !Dtran )
89         ferr( "Out of memory!!" );
90
91     for( i = Nstates ; --i >= 0 ; )
92     {
93         accept_states[i].string = Dstates[i].accept;
94         accept_states[i].anchor = Dstates[i].anchor;
95     }
96
97     free( Dstates );
98     *dfap      = Dtran;
99     *acceptp   = accept_states;
100
101    if( Verbose )
102    {
103        printf("\n%d out of %d DFA states in initial machine.\n",
104                           Nstates, DFA_MAX);
105
106        printf("%d bytes required for uncompressed tables.\n\n",
107               Nstates * MAX_CHARS * sizeof(TTYPE)                /* dtran */
108               + Nstates           * sizeof(TTYPE) );            /* accept */
109
110        if( Verbose > 1 )
111        {
112            printf("The un-minimized DFA looks like this:\n\n");
113            pheader(stdout, Dtran, Nstates, accept_states);
114        }
115    }
116

```



Listing 2.40. continued...

```
117     return Nstates ;
118 }
```

Several support functions are needed to do the work, all in Listing 2.40. The `add_to_dstates()` function on line 23 adds a new DFA state to the `Dstates` array and increments the number-of-states counter, `Nstates`. It returns the state number (the index in `Dstates`) of the newly added state. `in_dstates()` on line 31 is passed a set of NFA states, and returns the state number of an existing state that uses the same set (or -1 if there is no such state). The routine just does a linear search, which is probably not the best strategy here. Something like a binary tree would be better.

`add_to_dstates()`
`in_dstates()`

Listing 2.41. *dfa.c*—NFA to DFA Conversion: Support Functions

```
119 int add_to_dstates( NFA_set, accepting_string, anchor )
120 SET      *NFA_set;
121 char   *accepting_string;
122 int    anchor;
123 {
124     int nextstate;
125
126     if( Nstates > (DFA_MAX-1) )
127         ferr("Too many DFA states\n");
128
129     nextstate = Nstates++ ;
130     Dstates[ nextstate ].set      = NFA_set;
131     Dstates[ nextstate ].accept = accepting_string;
132     Dstates[ nextstate ].anchor = anchor;
133
134     return nextstate;
135 }
136
137 /*-----*/
138
139 PRIVATE int in_dstates( NFA_set )
140 SET      *NFA_set;
141 {
142     /* If there's a set in Dstates that is identical to NFA_set, return the
143      * index of the Dstate entry, else return -1.
144      */
145
146     DFA_STATE *p;
147
148     for( p = &Dstates[Nstates]; --p >= Dstates ; )
149         if( IS_EQUIVALENT( NFA_set, p->set ) )
150             return( p - Dstates );
151
152     return( -1 );
153 }
154
155 /*-----*/
156
157 PRIVATE DFA_STATE      *get_unmarked()
158 {
159     /* Return a pointer to an unmarked state in Dstates. If no such state
160      * exists, return NULL. Print an asterisk for each state to tell the
161      * user that the program hasn't died while the table is being constructed.
162      */

```



Listing 2.41. continued...

```

163
164     for(; Last_marked < &Dstates[Nstates] ; ++Last_marked )
165     {
166         if( !Last_marked->mark )
167         {
168             putc( '*' , stderr );
169             fflush( stderr );
170
171             if( Verbose > 1 )
172             {
173                 fputs("-----\n", stdout );
174                 printf("working on DFA state %d = NFA states: ",
175                         Last_marked-Dstates);
176                 pset( Last_marked->set, fprintf, stdout );
177                 putchar('\n');
178             }
179
180             return Last_marked;
181         }
182     }
183     return NULL;
184 }
185
186 /*----- */
187
188 PRIVATE void free_sets()
189 {
190     /* Free the memory used for the NFA sets in all Dstate entries. */
191
192     DFA_STATE *p;
193
194     for( p = &Dstates[Nstates]; --p >= Dstates ; )
195         delset( p->set );
196 }
```

get_unmarked()

A pointer to the next unmarked state is returned by `get_unmarked()`, declared on line 157. The `Last_marked` variable was initialized by `dfa()` on line 38 of Listing 2.40. Since new states are always added to the end of the table, everything above `Last_marked` will have been marked with a previous expansion, so these states don't have to be examined now. The asterisk that's printed on line 168 of Listing 2.41 actually has two purposes. It tells the user that the program hasn't hung (it takes a while to make the tables), and it lets the user terminate the program with a Ctrl-Break (a SIGINT under UNIX). It's needed because MS-DOS ignores Ctrl-Breaks until a system call of some sort is made by a program. No such call is made while creating the tables, but a system call is required to print the asterisks.

free_sets()

Finally, `free_sets()` on line 188 of Listing 2.41 goes through the entire `Dstates` array, deallocating the memory used for the sets of NFA states.

**NFA to DFA Conversion:
make_dtran().**

The actual NFA to DFA conversion is done by `make_dtran()`, in Listing 2.42. It is a straightforward implementation of the algorithm in Table 2.10 on page 124.

Listing 2.42. *dfa.c*—NFA to DFA Conversion: Conversion Function

```

197 PRIVATE void make_dtran( sstate )           /* Starting NFA state. */
198 int sstate;                                /* Set of NFA states that define */
199 {                                           /* the next DFA state. */
200     SET *NFA_set;                           /* State currently being expanded. */
201     DFA_STATE *current;                     /* Goto DFA state for current char. */
202     int nextstate;                         /* Current DFA state is an accept */
203     char *isaccept;                        /* (this is the accepting string). */
204     int anchor;                            /* Anchor point, if any. */
205     int c;                                 /* Current input character. */
206
207     /* Initially Dstates contains a single, unmarked, start state formed by
208      * taking the epsilon closure of the NFA start state. So, Dstates[0]
209      * (and Dtran[0]) is the DFA start state.
210
211     */
212
213     NFA_set = newset(); ;
214     ADD( NFA_set, sstate );
215
216     Nstates = 1;
217     Dstates[0].set = e_closure(NFA_set, &Dstates[0].accept, &Dstates[0].anchor);
218     Dstates[0].mark = 0;
219
220     while( current = get_unmarked() )          /* Make the table */
221     {
222         current->mark = 1;
223
224         for( c = MAX_CHARS ; --c >= 0 ; )
225         {
226             if( NFA_set = move(current->set, c) )
227                 NFA_set = e_closure( NFA_set, &isaccept, &anchor );
228
229             if( !NFA_set )                      /* no outgoing transitions */
230                 nextstate = F;
231
232             else if( (nextstate = in_dstates(NFA_set)) != -1 )
233                 delset( NFA_set );
234
235             else
236                 nextstate = add_to_dstates( NFA_set, isaccept, anchor );
237
238             Dtran[ current-Dstates ][ c ] = nextstate ;
239         }
240     }
241
242     putc( '\n', stderr );           /* Terminate string of '*'s printed in
243                                   * get_unmarked();
244                                   */
245
246     free_sets();      /* Free the memory used for the DFA_STATE sets */
247 }

```

2.5.7 DFA Minimization—Theory*

The DFA created earlier (in Figure 2.15 on page 123) is not an optimal DFA. Only six of the eight states are actually required. A somewhat simplified version of the transition matrix for the original machine is reproduced in Table 2.11. A quick inspection of this table shows that States 6 and 7 are equivalent because they have identical next-state transitions and they're both accepting states. If all transitions into State 6 were replaced with transitions to State 7, the machine would still accept the same strings as before. One additional state can be eliminated, however, and simple inspection is not good enough for finding this extra state. I'll discuss how to find the extra state in this section.

Table 2.11. DFA Transition Matrix for DFA in Figure 2.15

	Lookahead		Accepting
	D	.	
Current state	0	1	no
	1	4	no
	2	3	no
	3	—	yes
	4	4	no
	5	6	yes
	6	7	yes
	7	7	yes

Finding equivalent states.

You find equivalent states by systematically eliminating those states that can't be equivalent—partitioning the initial array into potentially equivalent states, and then gradually partitioning the partitions. When you're done, states that share a partition are equivalent. Initially, you partition the matrix into two parts, one for accepting states and another for nonaccepting states:

	D	.
0	1	2
1	4	5
2	3	—
4	4	2
3	—	—
5	6	—
6	7	—
7	7	—

Implicit failure state.

(The partition number is on the right.) The implicit failure state is alone in a special, unnumbered partition that's not shown in the foregoing table. The next step in the minimization process creates a set of new partitions by looking at the `goto` transitions for each state and eliminating states that are not equivalent. If the outgoing transitions from two states go to different partitions, then the states are not equivalent. You do the elimination on a column-by-column basis. Starting with the *D* column, States 0, 1, and 4 all go to a state in Partition 0 on a *D*, but State 2 goes to Partition 1 on a *D*. Consequently, State 2 must be removed from Partition 0. Formally, you say that State 4 is *distinguished* from the other states in the partition by a *D*. Continuing in this manner, State 3 is also distinguished from States 5, 6, and 7 by a *D* because State 3 goes to the failure state on a *D*, but the other states go to a state in Partition 1. The failure state is not in Partition 1, so State 3 is distinguished by a *D*.

...is distinguished...

D	.
0	1 2
1	4 5
4	4 2
2	3 -
3	- -
5	6 -
6	7 -
7	7 -

0

1

Now you go down the dot (.) column. The dot distinguishes State 1 from States 0 and 4 because State 1 goes to a state in Partition 1 on a dot, but States 0 and 4 both go to States in Partition 2. The new partitions are:

D	.
0	1 2
4	4 2
1	4 5
2	3 -
3	- -
5	6 -
6	7 -
7	7 -

0

4

2

3

4

5

1

Next, you go through the array a second time, column by column. Now, *D* distinguishes State 0 from State 4 because State 0 goes to a state in Partition 4 on a *D*, but State 4 goes to a state in Partition 0 on a *D*. No other states can be distinguished from each other, so we're done. The final partitions look like this:

D	.
0	1 2
4	4 2
1	4 5
2	3 -
3	- -
5	6 -
6	7 -
7	7 -

0

5

4

2

3

6

1

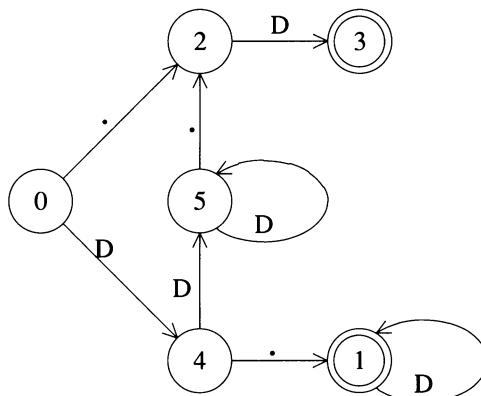
The next step is to build a new transition matrix. Each partition is a single state in the minimized DFA, and all next states in the original table are replaced by the number of the partition in which the state is found. For example, since States 5, 6, and 7 are all in Partition 1, all references to one of these states in the original table are replaced by a reference to the new State 1. The new table looks like this:

D		
0	4	2
1	1	-
2	3	-
3	-	-
4	5	1
5	5	2

The algorithm is formalized in Table 2.12. The transition diagram for the new state machine is in Figure 2.16.

Table 2.12. DFA-Minimization Algorithm

<i>c</i>	A character in the alphabet used by the DFA.
<i>group</i>	A set of potentially equivalent states (a partition of the original transition matrix).
<i>groups</i>	A collection of <i>groups</i> .
<i>new</i>	The set of states that have been distinguished from other states in the current group.
<i>first</i>	First state in a <i>group</i> .
<i>next</i>	One of the other states in a <i>group</i> , FALSE if there is no such state.
<i>goto_first</i>	A transition on <i>c</i> out of the <i>first</i> state comes here.
<i>goto_next</i>	A transition on <i>c</i> out of the <i>next</i> state comes here.
Initially:	Partition the original states into a series of <i>groups</i> . All nonaccepting states are in a single group, and accepting states that have the same accepting string are grouped together. A one-element group containing a single accepting state is permissible.
Repeat the following until no new groups are added to <i>groups</i> :	
<pre> for(each group in <i>groups</i>) { <i>new</i> = \emptyset <i>first</i> = the first state in the current group. <i>next</i> = the next state of the current group or FALSE if none. while(<i>next</i>) { for(each character <i>c</i>) { <i>goto_first</i> = state reached by making a transition on <i>c</i> out of <i>first</i>. <i>goto_next</i> = state reached by making a transition on <i>c</i> out of <i>next</i>. if(<i>goto_first</i> is not in the same group as <i>goto_next</i>) move <i>next</i> from the current group into <i>new</i> } <i>next</i> = the next state in the current group or FALSE if none. } if(<i>new</i> is not empty) add it to <i>groups</i> } </pre>	

Figure 2.16. Minimized DFA for $(D^* \setminus .DID \setminus D^*)$ 

2.5.8 DFA Minimization—Implementation

The algorithm in the previous section is implemented by the code starting in Listing 2.43. The `Groups[]` array on line 16 is the collection of partitions. The array index is used as the partition number, and each group is a SET of NFA state numbers. `Numgroups`, on line 17, is the number of groups currently in the array. `Ingroup[]`, on line 18, is the inverse of `Groups[]`. It is indexed by state number and evaluates to the group in which the state is found. All three of these are initialized to zeros by the access function `min_dfa()` (on lines 43 to 45). The DFA is created by the `dfa()` call on line 47, and it's minimized by the subroutine `minimize()`, called on the next line.

The `Groups[]` and `Ingroup[]` arrays are initialized in `init_groups()`, at the top of Listing 2.44. The program decides which group to place a state in by looking at the accepting-string pointer for that state. Since all nonaccepting states have NULL accepting strings, they'll all end up in a single group—Group 0 because DFA State 0 is always nonaccepting. The other groups in the initial partitioning contain accepting states. If two such states have the same accepting string, they are considered equivalent and are put into the same group. Note that since the starting state (0) always goes into Group 0, and since the minimization algorithm never moves the first element of the group to another partition, the starting state is still 0 after the minimization is complete.

The `goto` statement on line 78 of Listing 2.44 is necessary because it has to break out of a doubly nested loop in order to avoid forming a new group. The same thing could probably be done with a Boolean flag of some sort, but the `goto` seems cleaner.

The actual minimization is done by `minimize()` starting on line 104 of Listing 2.44. It is passed the number of states in the DFA transition table, a pointer to the DFA table itself, and a pointer to the accepting-state table. Note that there's an extra level of indirection in both of these last two parameters. That is, addresses of pointers to the two tables are passed to `minimize()`. I've done this because entirely new tables are created to replace the existing ones—`*dfap` and `*accept` are modified to point at the new tables. The memory used for the old tables is discarded automatically by `minimize()`.

The routine is a straightforward implementation of the algorithm in Table 2.12 on page 134. It terminates when an entire pass is made through the set of groups without adding a new group. The test on lines 143 to 146 is complicated by the failure transitions, which have no explicit existence. There is no physical group that holds the failure state, so I have to test here to make sure that neither of the two next states are failure

The partitions:
`Groups[], Ingroup[],`
`Numgroups`

Access function:
`min_dfa()`

Group initialization:
`init_groups()`.

Minimization: `minimize()`

Listing 2.43. *minimize.c*— Data Structures and Access Function

```

1 #include <stdio.h>
2 #ifdef MSDOS
3 #   include <stdlib.h>
4 #else
5 #   include <malloc.h>
6 #endif
7 #include <tools/debug.h>
8 #include <tools/set.h>
9 #include <tools/hash.h>
10 #include <tools/compiler.h>
11 #include "dfa.h"
12 #include "globals.h" /* externs for Verbose */
13
14 /* MINIMIZE.C: Make a minimal DFA by eliminating equivalent states. */
15
16 PRIVATE SET *Groups [ DFA_MAX ]; /* Groups of equivalent states in Dtran */
17 PRIVATE int Numgroups; /* Number of groups in Groups */
18 PRIVATE int Ingroup [ DFA_MAX ]; /* the Inverse of Group */
19
20 /-----
21 * Prototypes for subroutines in this file:
22 */
23
24 PRIVATE void fix_dtran ( ROW*[ ], ACCEPT** );
25 PRIVATE void init_groups ( int, ACCEPT* );
26 PUBLIC int min_dfa ( char * (*) (), ROW*[ ], ACCEPT** );
27 PRIVATE void minimize ( int, ROW*[ ], ACCEPT** );
28 PRIVATE void pgroups ( int );
29
30 /-----
31
32 PUBLIC int min_dfa( ifunct, dfap, acceptp )
33     *( *ifunct )();
34 ROW *( dfap[] );
35 ACCEPT *( *acceptp );
36 {
37     /* Make a minimal DFA, eliminating equivalent states. Return the number of
38     * states in the minimized machine. *sstatep = the new start state.
39     */
40     int nstates; /* Number of DFA states */
41
42     memset( Groups, 0, sizeof(Groups) );
43     memset( Ingroup, 0, sizeof(Ingroup) );
44     Numgroups = 0;
45
46     nstates = dfa( ifunct, dfap, acceptp );
47     minimize ( nstates, dfap, acceptp );
48
49     return Numgroups;
50 }
51 }
```

transitions before allowing an index into the `Ingroup[]` array.

Removing redundant
states: `fix_dtran()`

The `fix_dtran()` routine on line 177 of Listing 2.44 uses the previously-computed `Groups[]` to remove redundant states from the table. It's difficult to do the compression in place because the original states are randomly distributed throughout the groups, so two new arrays are allocated on lines 199 and 200. The original arrays are

freed on lines 218 and 219, and the array pointers are changed to point at the new arrays on the next couple of lines.

The remainder of *minimize.c* consists of a few debugging routines in Listing 2.45. *pgroups()* prints out all the existing groups and is used for very-verbose-mode (-V) diagnostics.

Print the groups:
pgroups()

Listing 2.44. *minimize.c*—Minimization Functions

```

52 PRIVATE void init_groups( nstates, accept )
53 ACCEPT *accept;
54 {
55     SET **last ;
56     int i, j ;
57
58     last      = & Groups[0] ;
59     Numgroups = 0;
60
61     for( i = 0; i < nstates ; i++ )
62     {
63         for( j = i; --j >= 0 ;)
64         {
65             /* Check to see if a group already exists that has the same
66              * accepting string as the current state. If so, add the current
67              * state to the already existing group and skip past the code that
68              * would create a new group. Note that since all nonaccepting states
69              * have NULL accept strings, this loop puts all of these together
70              * into a single group. Also note that the test in the for loop
71              * always fails for group 0, which can't be an accepting state.
72             */
73
74             if( accept[i].string == accept[j].string )
75             {
76                 ADD( Groups[ Ingroup[j] ], i );
77                 Ingroup[i] = Ingroup[j];
78                 goto match;
79             }
80         }
81
82         /* Create a new group and put the current state into it. Note that ADD()
83          * has side effects, so "last" can't be incremented in the ADD
84          * invocation.
85         */
86
87         *last = newset();
88         ADD( *last, i );
89         Ingroup[i] = Numgroups++;
90         ++last;
91
92     match:; /* Group already exists, keep going */
93     }
94
95     if( Verbose > 1 )
96     {
97         printf ( "Initial groupings:\n" );
98         pgroups ( nstates );
99     }
100 }
```



Listing 2.44. continued...

```

102  /*-----*/
103
104 PRIVATE void minimize( nstates, dfap, acceptp )
105 int      nstates;           /* number of states in dtran[] */      */
106 ROW     *( dfap[] ) ;       /* DFA transition table to compress */   */
107 ACCEPT  *( *acceptp );     /* Set of accept states */                  */
108 {
109     int old_numgroups; /* Used to see if we did anything in this pass. */
110     int c;               /* Current character. */
111     SET **current;      /* Current group being processed. */
112     SET **new ;         /* New partition being created. */
113     int first;          /* State # of first element of current group. */
114     int next;           /* State # of next element of current group. */
115     int goto_first;     /* target of transition from first[c]. */
116     int goto_next;      /* other target of transition from first[c]. */
117
118     ROW    *dtran = *dfap;
119     ACCEPT *accept = *acceptp;
120
121     init_groups( nstates, accept );
122     do
123     {
124         old_numgroups = Numgroups;
125         for( current = &Groups[0]; current < &Groups[Numgroups]; ++current )
126         {
127             if( num_ele( *current ) <= 1 )
128                 continue;
129
130             new = &Groups[ Numgroups ];
131             *new = newset();
132
133             next_member( NULL );
134             first = next_member( *current );
135
136             while( (next = next_member(*current)) >= 0 )
137             {
138                 for( c = MAX_CHARS; --c >= 0 ; )
139                 {
140                     goto_first = dtran[ first ][ c ];
141                     goto_next = dtran[ next ][ c ];
142
143                     if( goto_first != goto_next
144                         &&( goto_first == 'F'
145                             || goto_next == F
146                             || Ingroup[goto_first] != Ingroup[goto_next]
147                             )
148                     )
149                     {
150                         REMOVE( *current, next ); /* Move the state to */
151                         ADD   ( *new,      next ); /* the new partition */
152                         Ingroup[ next ] = Numgroups ;
153                         break;
154                     }
155                 }
156             }
157
158             if( IS_EMPTY( *new ) )
159                 delset( *new );

```



Listing 2.44. continued...

```

160         else
161             ++Numgroups;
162     }
163
164 } while( old_numgroups != Numgroups );
165
166 if( Verbose > 1 )
167 {
168     printf("\nStates grouped as follows after minimization:\n");
169     pgroups( nstates );
170 }
171
172 fix_dtran( dfap, acceptp );
173 }
174
175 /*-----*/
176
177 PRIVATE void fix_dtran( dfap, acceptp )
178 ROW    *( dfap[] );
179 ACCEPT  *( *acceptp );
180 {
181     /* Reduce the size of the dtran, using the group set made by minimize().
182      * Return the state number of the start state. The original dtran and accept
183      * arrays are destroyed and replaced with the smaller versions.
184      * Consider the first element of each group (current) to be a
185      * "representative" state. Copy that state to the new transition table,
186      * modifying all the transitions in the representative state so that they'll
187      * go to the group within which the old state is found.
188 */
189
190     SET    **current    ;
191     ROW    *newdtran   ;
192     ACCEPT *newaccept  ;
193     int    state       ;
194     int    i           ;
195     int    *src, *dest ;
196     ROW    *dtran     = *dfap;
197     ACCEPT *accept   = *acceptp;
198
199     newdtran = (ROW    *) calloc( Numgroups, sizeof(ROW    ) );
200     newaccept = (ACCEPT *) calloc( Numgroups, sizeof(ACCEPT) );
201
202     if( !newdtran || !newaccept )
203         ferr("Out of memory!!!!");
204
205     next_member( NULL );
206     for( current = &Groups[Numgroups]; --current >= Groups; )
207     {
208         dest   = &newdtran[ current-Groups ][ 0 ];
209         state  = next_member(*current);          /* All groups have at */
210         src    = & dtran[ state ][ 0 ];           /* least one element. */
211
212         newaccept[ current-Groups ]=accept[state]; /* Struct. assignment */
213
214         for( i = MAX_CHARS; --i >= 0 ; src++, dest++ )
215             *dest = (*src == F) ? F : Ingroup[ *src ];
216     }
217

```



Listing 2.44. continued...

```

218     free( *dfap    );
219     free( *acceptp );
220     *dfap    = newdtran ;
221     *acceptp = newaccept ;
222 }
```

Listing 2.45. minimize.c— Debugging Routines

```

223 PRIVATE void pgroups( nstates )
224 int      nstates ;
225 {
226     /* Print all the groups used for minimization. */
227
228     SET **current;
229
230     for( current = &Groups[Numgroups] ; --current >= Groups ; )
231     {
232         printf( "\tgroup %d: {", current - Groups );
233         pset( *current, fprintf, stdout );
234         printf( "}\n" );
235     }
236     printf("\n");
237     while( --nstates >= 0 )
238         printf("\tstate %2d is in group %2d\n", nstates, Ingroup[nstates] );
239 }
```

Dead states.

Unreachable states.

There are actually two additional minimizations that could be done here and are not: *dead states*—nonaccepting states that have no outgoing transitions other than ones that loop back into the current state—can be removed, as can *unreachable states*—states with no incoming edges. In practice, both types of states are rare, so it didn't seem worth the effort to remove them. I'll leave this further minimization as an exercise.

2.5.9 Compressing and Printing the Tables

The remainder of this chapter describes the remainder of L^EX—primarily housekeeping and table-printing functions. Skip to the next chapter if you're not following the implementation details.

This section discusses three table-output routines: one for uncompressed tables, one for pair-compressed tables, and a third routine that eliminates redundant rows and columns and prints the result. The first two of these are of sufficient generality to be put in a library. The third routine is used only by L^EX.

Generate yy_next()
function: defnext()

Printing two-dimensional
arrays: print_array()

2.5.9.1 Uncompressed Tables. The uncompressed tables are printed by two subroutines in two separate input files. The defnext() routine in Listing 2.46 prints the yy_next() macro definition—yy_next() figures the next state from the current state and input character. The array access is hidden in a macro so that the state-machine driver can call yy_next() for the next state, regardless of the compression method used. Each compression method supplies a different version of yy_next().

The uncompressed two-dimensional array is output by print_array() in Listing 2.47. The routine is passed a FILE pointer for the output stream, a pointer to the array, and the two dimensions. It is configured to print an array of ints, but you can change

Listing 2.46. *defnext.c*—Print Next-State Macro: Uncompressed Table

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/compiler.h>      /* needed only for prototype */
4
5 PUBLIC void defnext(fp, name)
6 FILE *fp;
7 char *name;
8 {
9     /* Print the default yy_next(s,c) subroutine for an uncompressed table. */
10
11     static char *comment_text[] =
12     {
13         "yy_next(state,c) is given the current state and input character and",
14         "evaluates to the next state.",
15         NULL
16     };
17
18     comment( fp, comment_text );
19     fprintf( fp, "#define yy_next(state, c) %s[ state ][ c ]\n", name );
20 }
```

this type by modifying the **typedef** for ATYPE on line eight. The array is treated internally as a one-dimensional array having nrows × ncols cells. The main use of the row size (nrows) is to print the inner brackets on lines 28 and 43. NCOLS, defined on line nine, determines how many array elements are printed on each output line. It controls formatting only—not any of the actual array dimensions. `print_array()` just prints the initialization part of the declaration (and the terminating semicolon). It doesn't print the actual declaration. Use it like this:

```
printf("unsigned char Array[%d][%d]=\n", nrows, ncols );
print_array( stdout, array, nrows, ncols );
```

The declared type of the output array (**unsigned char** above) is immaterial, even though the input array must be **ints**. It's up to you to make sure that all values in the input array fit into the cells of the output array. In the current example, no cell in the input array can have a value greater than 255 or the number will effectively be truncated.

2.5.9.2 Pair-Compressed Tables. The next set of routines, in Listing 2.48, produce pair-compressed tables, as were pictured in Figure 2.7 on page 70. Various macros on lines 21 to 26 control the names and types used for both the input and output array. ATYPE and NCOLS are used as before: they are the declared type of the input array and the number of columns on one line of the output array. The other three definitions are used for the storage class of the array and the type of one array cell. The printed tables assume that the table type and storage class are defined somewhere—you have to put something like the following at the top of the file in which the tables are placed:

```
typedef unsigned char YY_TTYPE;
#define YYPRIVATE static
```

Note that the `yy_next()`, which figures the next state, can be declared with a different storage class than the tables themselves (though it isn't in the current configuration). This way you can have a public decoding routine and **static** tables.

The table-generation routine, `pairs()`, starts on line 29 of Listing 2.48. It is passed three more arguments than `print_array()`. The name argument is the name used for the output arrays. The names for the arrays that represent rows are manufactured by

ATYPE

NCOLS

ATYPE, NCOLS

Printing pair-compressed
tables: `pairs()`

Listing 2.47. print_ar.c—Print A Two Dimensional Array

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/compiler.h>           /* for prototypes only */
4 /-----
5 * PRINT_AR.C: General-purpose subroutine to print out a 2-dimensional array.
6 */
7
8 typedef int      ATYPE;
9 #define NCOLS    10   /* Number of columns used to print arrays      */
10 /*-----*/
11
12 PUBLIC void     print_array( fp, array, nrows, ncols )
13 FILE      *fp;
14 ATYPE     *array;          /* DFA transition table                      */
15 int       nrows;         /* Number of rows    in array[]            */
16 int       ncols;         /* Number of columns in array[]           */
17 {
18     /* Print the C source code to initialize the two-dimensional array pointed
19      * to by "array." Print only the initialization part of the declaration.
20      */
21
22     int       i;
23     int       col;        /* Output column           */
24
25     fprintf( fp, "{\n" );
26     for( i = 0; i < nrows ; i++ )
27     {
28         fprintf(fp, /* %02d */ { , i );
29
30         for( col = 0; col < ncols; col++ )
31         {
32             fprintf(fp, "%3d" , *array++ );
33             if( col < ncols-1 )
34                 fprintf(fp, ", " );
35
36             if( (col % NCOLS) == NCOLS-1 && col != ncols-1 )
37                 fprintf(fp, "\n" );
38         }
39
40         if( col > NCOLS )
41             fprintf( fp, "\n" );
42
43         fprintf( fp, " }%c\n", i < nrows-1 ? ',' : ' ' );
44     }
45     fprintf(fp, "};\n");
46 }
```

threshold argument to
pairs()

putting a numeric suffix at the end of this string. The array of pointers to rows uses the string without modification. The threshold argument determines the point at which character/next-state pairs are abandoned in favor of a one-dimensional array. If the number of transitions in a row are less than or equal to threshold, then pairs are used, otherwise an array is used. Finally, numbers controls the output format of the character part of the pair. If it's true, a decimal number is used (an ASCII 'a', for example, is output as 97); if false, character constants are used (an 'a' is printed as an 'a').

The pairs() routine returns the number of cells used to store the row arrays. The number of bytes used by the entire table is:

$$(number_of_cells \times \text{sizeof}(\text{TYPE})) + (\text{nrows} \times \text{sizeof}(\text{TYPE}^*))$$

The rows are printed in the loop that starts on line 52 of Listing 2.48. The loop on line nine counts the number of nonfailure transitions in the row, and the row is generated only if there are nonfailure transitions. The elements themselves are printed in the loop on line 75. The second half of the subroutine, on lines 110 to 129 prints out the array of pointers to the rows. The pnext() subroutine that starts on line 135 prints the next-state subroutine (it's too big to be a macro).

Listing 2.48. pairs.c—Pair Compression

```

1  #include <stdio.h>
2  #include <tools/debug.h>
3  #include <tools/compiler.h>           /* for prototypes only */
4
5  /* PAIRS.C      This module contains the routines to compress a table
6   *               horizontally (using char/next pairs) and then print the
7   *               compressed table. The compressed array looks like this:
8   *
9   *   Yy_nxt:      Yy_nxtDD:
10  *   +-----+      +-----+
11  *   |  *--|--->|  0   | Next state array, indexed by character |
12  *   +-----+      +-----+
13  *   |          |
14  *   +-----+      +-----+-----+-----+-----+-----+
15  *   |  *--|--->| count | c1 | s1 | c2 | s2 | ... |
16  *   +-----+      +-----+-----+-----+-----+-----+
17  *   | NULL |          (there are no non-failure transitions if NULL)
18  *   +-----+
19  */
20
21  typedef int      ATYPE;      /* Declared type of input tables        */
22
23  #define NCOLS     10          /* Number of columns used to print arrays */
24  #define TYPE       "YY_TTYPE"    /* Declared type of output tables.      */
25  #define SCLASS     "YYPPRIVATE"  /* Storage class of all the tables.    */
26  #define D_SCLASS   "YYPPRIVATE" /* Storage class of the decoding routine */
27  /*-----*/
28
29  PUBLIC int pairs( fp, array, nrows, ncols, name, threshold, numbers )
30
31  FILE    *fp;                /* output file                         */
32  ATYPE   *array;             /* DFA transition table                 */
33  int     nrows;              /* Number of rows in array[]          */
34  int     ncols;              /* Number of columns in array[]        */
35  char    *name;              /* Name used for output array.        */
36  int     threshold;          /* Array vs. pairs threshold         */
37  int     numbers;            /* Use numbers for char. part of pair */
38  {
39      /* Generate the C source code for a pair-compressed DTRAN. Returns the
40      * number of cells used for the YysDD arrays. The "numbers" argument
41      * determines the output format of the character part of a
42      * character/next-state pair. If numbers is true, then normal numbers
43      * are used, otherwise ASCII characters are used. For example: 'a',100
44      * as compared to: 97,100 ('a' == 0x61 == 97, decimal).
45      */
46
47  int     i, j, ntransitions, nprinted, ncommas;
48  int     num_cells = 0;          /* cells used for rows */
49  ATYPE   *p;

```



Listing 2.48. continued...

```

50     char    *bin_to_ascii();
51
52     for( i = 0; i < nrows ; i++ )
53     {
54         ntransitions = 0;
55         for( p = array + (i * ncols), j = ncols; --j >= 0; p++ )
56             if( *p != -1 )
57                 ++ntransitions ;
58
59         if( ntransitions > 0 )
60         {
61             fprintf(fp, "%s %s %s%-2d[] = { ", SCLASS, TYPE, name, i ); /*}*/
62             ++num_cells;
63             if( ntransitions > threshold ) /* array */
64                 fprintf(fp, "0,\n");
65             else /* pairs */
66             {
67                 fprintf(fp, "%2d, ", ntransitions );
68                 if( threshold > 5 )
69                     fprintf(fp, "\n" );
70             }
71
72         nprinted = NCOLS;
73         ncommas = ntransitions;
74
75         for( p = array + (i * ncols), j = 0; j < ncols; j++, p++ )
76         {
77             if( ntransitions > threshold ) /* array */
78             {
79                 ++num_cells ;
80                 --nprinted ;
81
82                 fprintf(fp, "%3d" , *p );
83                 if( j < ncols-1 )
84                     fprintf(fp, ", " );
85             }
86             else if( *p != -1 ) /* pairs */
87             {
88                 num_cells += 2;
89
90                 if( numbers )
91                     fprintf(fp, "%d,%d", j, *p );
92                 else
93                     fprintf(fp, "'%s',%d", bin_to_ascii(j,0), *p );
94
95                 nprinted -= 2;
96                 if( --ncommas > 0 )
97                     fprintf(fp, ", " );
98             }
99
100            if( nprinted <= 0 )
101            {
102                fprintf(fp, "\n");
103                nprinted = NCOLS;
104            }
105        }
106        fprintf(fp, "};\n" );
107    }
108 }
109 }
```



Listing 2.48. continued...

```

110     fprintf(fp, "\n%s %s *%s[ %d ] =\n{\n      ", SCLASS, TYPE, name, nrows );
111
112     nprinted = 10;                                /* 10 elements on a line */
113     for( --nrows, i = 0; i < nrows ; i++ )
114     {
115         ntransitions = 0;
116         for( p = array + (i * ncols), j = ncols; --j >= 0; p++ )
117             if( *p != -1 )
118                 ++ntransitions ;
119
120         fprintf(fp, ntransitions ? "%s%-2d, " : "NULL,  ", name, i );
121
122         if( --nprinted <= 0 )
123         {
124             fprintf(fp, "\n      ");
125             nprinted = 10;
126         }
127     }
128
129     fprintf(fp, "%s%-2d\n};\n\n", name, i );
130     return num_cells;
131 }
132
133 /-----*/
134
135 PUBLIC void pnext( fp, name )
136 char *name;
137 FILE *fp;
138 {
139     /* Print out a next(state,c) subroutine for a table compressed      */
140     /* into char/next-state pairs by the routines in pairs.c.          */
141
142     static char *toptext[]=
143     {
144         "unsigned int    c ;",
145         "int           cur_state ;",
146         "{",
147         "    /* Given the current state and the current input character, return",
148         "     * the next state.",
149         "    */",
150         "",
151         NULL
152     };
153     static char *boptext[]=
154     {
155         "    int i;",
156         "",
157         "    if( p )",
158         "    {",
159         "        if( (i = *p++) == 0 )",
160         "            return p[ c ];",
161         "",
162         "        for( ; --i >= 0 ; p += 2 )",
163         "            if( c == p[0] )",
164         "                return p[1];",
165         "    }",
166         "    return YYF;",
167         "}",
168         NULL
169     };

```



Listing 2.48. continued...

```

170     fprintf (fp, "\n/*-----*/\n");
171     fprintf (fp, "%s %s yy_next( cur_state, c )\n", D_SCLASS, TYPE );
172     printv (fp, toptext );
173     fprintf (fp, "    %s *p = %s[ cur_state ] ;\n", TYPE, name );
174     printv (fp, boptext );
175 }
```

2.5.9.3 Redundant-Row-and-Column-Compressed Tables. The routines to print the array with redundant row and columns removed are in Listing 2.49. Since these are used only in LEX, I've taken less pains to make them as flexible as possible. The subroutine goes through the array column by column using the algorithm in Table 2.13. The same algorithm is used to compress the rows. Note that the row compression creates a single row that contains nothing but failure transitions. This approach seems better than setting the map-array entries to a failure marker because it makes the array access both more consistent and faster at the cost of only a few bytes—you don't have to test for failure before indexing into the map array.

Table 2.13. Algorithm to Eliminate Redundant Columns

<i>array</i>	array to compress
<i>map</i>	the map array being created
<i>save</i>	is the set of columns that are not redundant
<i>i, j</i>	counters
Initially:	<i>save</i> is empty.
	<pre> for(i = each column in array) { while(i is a column that's already been marked as redundant) ++i; add i to save set; j = i+1 ; while(j < rightmost column) if(column i ≡ column j) { column map[j] = i; mark j as redundant; } } j = 0; for(i = each element of save set) { copy column i to column j; ++j; } }</pre>

Listing 2.49. *squash.c*—Redundant Row and Column Elimination

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #ifndef MSDOS
4 # include <stdlib.h>
5 #else
6 # include <malloc.h>
7 #endif
8 #include <tools/debug.h>
9 #include <tools/set.h>
10 #include <tools/compiler.h>
11
12 #include "dfa.h"
13 #include "globals.h"
14
15 /* SQUASH.C -- This module contains the routines to compress a table
16 * horizontally and vertically by removing redundant columns and rows, and then
17 * print the compressed table. I haven't been as careful about making this
18 * routine general purpose because it's only useful to LeX. The pairs
19 * compression routines in pairs.c are used to compress the occs- and
20 * llama-generated tables so they're a little more complicated.
21 */
22
23 PRIVATE int Col_map[ MAX_CHARS ];
24 PRIVATE int Row_map[ DFA_MAX ];
25
26 #define NCOLS 16
27 #define TYPE "YY_TTYPE" /* Declared type of output tables. */ */
28 #define SCLASS "YYPRIVATE" /* Storage class of all the tables */ */
29
30 /-----
31 * Subroutine in this file:
32 */
33
34 PRIVATE int col_equiv ( int*, int*, int );
35 PRIVATE void col_cpy ( int*, int*, int , int , int );
36 PRIVATE void reduce ( ROW*, int*, int* );
37 PRIVATE void print_col_map ( FILE* );
38 PRIVATE void print_row_map ( FILE*, int );
39 PUBLIC void pmap ( FILE*, int*, int );
40 PUBLIC void cnext ( FILE*, char* );
41 PUBLIC int squash ( FILE*, ROW*, int , int , char* );
42
43 /-----*/
44
45 #define ROW_EQUIV(r1,r2,ncols) (memcmp( r1, r2, ncols * sizeof(int)) == 0 )
46 #define ROW_CPY(r1,r2,ncols) (memcpy( r1, r2, ncols * sizeof(int)) )
47
48 /-----*/
49
50 PUBLIC int squash( fp, dtran, nrows, ncols, name )
51 FILE *fp;
52 ROW *dtran;
53 char *name;
54 {
55     /* Compress (and output) dtran using equivalent-column elimination.
56     * Return the number of bytes required for the compressed tables
57     * (including the map but not the accepting array).
58     */
59

```



Listing 2.49. continued...

```

60     int oncols = ncols;                                /* original column count */
61     int onrows = nrows;                               /* original row    count */
62
63     reduce( dtran, &nrows, &ncols );                  /* Compress the tables   */
64
65     print_col_map ( fp );
66     print_row_map ( fp, onrows );
67
68     fprintf(fp, "%s %s %s[ %d ][ %d ]=\n", SCLASS, TYPE, name, nrows, ncols);
69     print_array( fp, (int *)dtran, nrows, ncols );
70
71     return( ( nrows * ncols * sizeof(TTYPE))           /* dtran   */
72             +(          onrows * sizeof(TTYPE))        /* row map */
73             +(          oncols * sizeof(TTYPE)) );      /* col map */
74 }
75
76 /*-----*/
77
78 PRIVATE int col_equiv( col1, col2, nrows )
79 int *col1, *col2;
80 {
81     /* Return 1 if the two columns are equivalent, else return 0 */
82
83     while( --nrows >= 0 && *col1 == *col2 )
84     {
85         col1 += MAX_CHARS;           /* Advance to next cell in the column */
86         col2 += MAX_CHARS;
87     }
88
89     return( !(nrows >= 0) );
90 }
91
92 /*-----*/
93
94 PRIVATE void col_cpy( dest, src, nrows, n_src_cols, n_dest_cols )
95 int *dest;                      /* Top of destination column           */
96 int *src;                       /* Top of source column               */
97 int nrows;                      /* Number of rows                   */
98 int n_src_cols;                 /* Number of columns in source array */
99 int n_dest_cols;                /* Number of columns in destination array */
100 {
101     /* Copy a column from one array to another. Both arrays are nrows deep,
102      * the source array is n_src_cols wide and the destination array is
103      * n_dest_cols wide.
104 */
105
106     while( --nrows >= 0 )
107     {
108         *dest = *src;
109         dest += n_dest_cols;
110         src += n_src_cols;
111     }
112 }
113
114 /*-----*/
115
116 PRIVATE void    reduce( dtran, p_nrows, p_ncols )
117 ROW     *dtran;                  /* DFA transition table               */
118 int      *p_nrows;                /* # of states in dtran              */
119 int      *p_ncols;                /* Pointer to column count           */

```



Listing 2.49. continued...

```

120  {
121      /* Reduce dtran horizontally and vertically, filling the two map arrays
122      * with the character mapping to the final, reduced transition table.
123      * Return the number of columns in the reduced dtran table.
124      *
125      * Col_map is the x (character) axis, Row_map is the y (next state) axis.
126      */
127
128      int ncols = *p_ncols;    /* number of columns in original machine */
129      int nrows = *p_nrows;    /* number of rows   in original machine */
130      int r_ncols;           /* number of columns in reduced machine */
131      int r_nrows;           /* number of rows   in reduced machine */
132      SET *save;             /* rows or columns that will remain in table */
133      int *current;          /* first of several identical columns */
134      int *compressed;        /* pointer to compressed array */
135      int *p;
136      int i, j;
137
138      /* ..... .
139      * First do the columns.
140      */
141
142      memset( Col_map, -1, sizeof(Col_map) );
143      save = newset();
144
145      for( r_ncols = 0 ;; r_ncols++ )
146      {
147          /* Skip past any states in the Col_map that have already been
148          * processed. If the entire Col_map has been processed, break.
149          */
150
151          for(i = r_ncols; Col_map[i] != -1 && i < ncols ; i++ )
152              ;
153
154          if( i >= ncols )
155              break;
156
157          /* Add the current column to the save set. It eventually ends up
158          * in the reduced array as column "r_ncols" so modify the Col_map
159          * entry accordingly. Now, scan through the array looking for
160          * duplicates of the current column (pointed to by current). If you
161          * find a duplicate, make the associated Col_map entry also point to
162          * "r_ncols."
163          */
164
165          ADD( save, i );
166          Col_map[i] = r_ncols;
167
168          current = &dtran[0][i];
169          p       = current + 1;
170
171          for( j = i; ++j < ncols ; p++ )
172              if( Col_map[j]==-1 && col_equiv(current, p, nrows) )
173                  Col_map[j] = r_ncols ;
174      }
175
176      /* Compress the array horizontally by removing all of the columns that
177      * aren't in the save set. We're doing this by moving all the columns
178      * that are in the save set to the proper position in a newly allocated
179      * array. You can't do it in place because there's no guarantee that the

```

Listing 2.49. continued...

```

180     * equivalent rows are next to each other.
181     */
182
183     if( !(compressed = (int *) malloc(nrows * r_ncols * sizeof(int))) )
184         ferr("Out of memory");
185
186     p = compressed;
187     for( next_member(NULL); (i = next_member(save)) != -1; )
188         col_cpy( p++, &dtran[0][i], nrows, ncols, r_ncols );
189
190     /* .....
191      * Eliminate equivalent rows, working on the reduced array
192      * created in the previous step. The algorithm used is the
193      * same.
194     */
195
196     memset( Row_map, -1, sizeof(Row_map) );
197     CLEAR( save );
198
199     for( r_nrows = 0 ;; r_nrows++ )
200     {
201         for( i = r_nrows; Row_map[i] != -1 && i < nrows; i++ )
202             ;
203
204         if( i >= nrows )
205             break;
206
207         ADD( save, i );
208         Row_map[i] = r_nrows;
209
210         current = compressed + ((i ) * r_ncols );
211         p       = compressed + ((i+1) * r_ncols );
212
213         for( j = i; ++j < nrows ; p += r_ncols )
214             if( Row_map[j]==-1 && ROW_EQUIV(current, p, r_ncols) )
215                 Row_map[j] = r_nrows ;
216     }
217
218     /*
219      * Actually compress rows, copying back into the original array space.
220      * Note that both dimensions of the array have been changed.
221     */
222
223     p = (int *)dtran;
224
225     for( next_member(NULL); (i = next_member(save)) != -1; )
226     {
227         ROW_CPY( p, compressed + (i * r_ncols) , r_ncols );
228         p += r_ncols;
229     }
230
231     delset ( save );
232     free   ( compressed );
233
234     *p_ncols = r_ncols ;
235     *p_nrows = r_nrows ;
236 }
237

```



Listing 2.49. continued...

```

238 /*-----*/
239
240 PRIVATE void print_col_map( fp )
241 FILE *fp;
242 {
243     static char *text[] =
244     {
245         "The Yy_cmap[] and Yy_rmap arrays are used as follows:",
246         "",
247         " next_state= Yydtran[ Yy_rmap[current_state] ][ Yy_cmap[input_char] ];",
248         "",
249         "Character positions in the Yy_cmap array are:",
250         "",
251         " ^@ ^A ^B ^C ^D ^E ^F ^G ^H ^I ^J ^K ^L ^M ^N ^O",
252         " ^P ^Q ^R ^S ^T ^U ^V ^W ^X ^Y ^Z ^[ ^\\ ^] ^^ ^`",
253         " ! \\" # $ % & ' ( ) * + , - . /",
254         " 0 1 2 3 4 5 6 7 8 9 : ; < = > ?",
255         " @ A B C D E F G H I J K L M N O",
256         " P Q R S T U V W X Y Z [ \\ ] ^",
257         " ` a b c d e f g h i j k l m n o",
258         " p q r s t u v w x y z { | } ~ DEL",
259         NULL
260     };
261
262     comment(fp, text);
263     fprintf(fp, "%s %s Yy_cmap[%d] =\n\n", SCLASS, TYPE, MAX_CHARS );
264     pmap( fp, Col_map, MAX_CHARS );
265 }
266
267
268 PRIVATE void print_row_map( fp, nrows )
269 FILE *fp;
270 int nrows;
271 {
272     fprintf(fp, "%s %s Yy_rmap[%d] =\n\n", SCLASS, TYPE, nrows );
273     pmap( fp, Row_map, nrows );
274 }
275
276 PRIVATE void pmap( fp, p, n )
277 FILE *fp; /* output stream */
278 int *p; /* pointer to array */
279 int n; /* array size */
280 {
281     /* Print a one-dimensional array.
282     */
283
284     int j;
285     for( j = 0; j < (n - 1); j++ )
286     {
287         fprintf(fp, "%3d," , *p++ );
288
289         if( (j % NCOLS) == NCOLS-1 )
290             fprintf(fp, "\n      ");
291     }
292
293     fprintf( fp, "%3d\n};\n\n", *p );
294 }
295

```



Listing 2.49. continued...

```

296  /*-----*/
297
298 PUBLIC void cnext( fp, name )
299 FILE *fp;
300 char *name;
301 {
302     /* Print out a yy_next(state,c) subroutine for the compressed table.
303      */
304
305     static char *text[] =
306     {
307         "yy_next(state,c) is given the current state number and input",
308         "character and evaluates to the next state.",
309         NULL
310     };
311
312     comment( fp, text );
313     fprintf( fp,
314             "#define yy_next(state,c) (%s[ Yy_xmap[state] ][ Yy_cmap[c] ])\\n",
315             name
316     );
317 }
```

2.5.10 Tying It All Together

LEX is completed by the subroutines presented in this section. (in Listings 2.50, 2.51, 2.52, and 2.53). These listings are commented sufficiently that additional comments here would be tedious. Read the code if you're interested.

I should also add, at this point, that there are more efficient methods than the one I've just described for converting a regular expression to a finite automata. (See [Aho] pp. 134–141.) I've chosen the current method because it serves to introduce several concepts (like closure) that will be useful later on.

Listing 2.50. signon.c—Print Sign-on Message

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3
4
5 signon()
6 {
7     /* Print the sign-on message. Since the console is opened explicitly, the
8      * message is printed even if both stdout and stderr are redirected.
9      */
10
11     FILE *screen;
12
13     UX( if( !(screen = fopen("/dev/tty", "w")) ) )
14     MS( if( !(screen = fopen("con:", "w")) ) )
15         screen = stderr;
16
17     /* The ANSI __DATE__ macro yields a string of the form: "Sep 01 1989". */
18     /* The __DATE__+7 gets the year portion of that string. */
19
```



Listing 2.50. continued...

```

20     fprintf(screen,"LeX 1.0 [%s]. (c) %s, Allen I. Holub. "
21                     "All rights reserved.\n", __DATE__, __DATE__+7 );
22
23     if( screen != stderr )
24         fclose(screen);
25 }
```

Listing 2.51. print.c— Print Remainder of Output File

```

1  #include <stdio.h>
2  #include <ctype.h>
3  #include <tools/debug.h>
4  #include <tools/set.h>
5  #include <tools/compiler.h>
6  #include "dfa.h"
7  #include "nfa.h"
8  #include "globals.h"
9
10 /* PRINT.C: This module contains miscellaneous print routines that do
11 * everything except print the actual tables.
12 */
13
14 PUBLIC void pheader( FILE *fp,      ROW dtran[], int nrows, ACCEPT *accept);
15 PUBLIC void pdriver( FILE *output,   int nrows, ACCEPT *accept);
16 /*-----*/
17
18 PUBLIC void pheader( fp, dtran, nrows, accept )
19 FILE    *fp;                      /* output stream */                      */
20 ROW     dtran[];                 /* DFA transition table */                  */
21 int    nrows;                   /* Number of states in dtran[] */          */
22 ACCEPT *accept;                /* Set of accept states in dtran[] */        */
23 {
24     /* Print out a header comment that describes the uncompressed DFA. */
25
26     int      i, j;
27     int      last_transition ;
28     int      chars_printed;
29     char    *bin_to_ascii() ;
30
31     fprintf(fp, "#ifdef __NEVER__\n" );
32     fprintf(fp, "/*-----\n");
33     fprintf(fp, " * DFA (start state is 0) is:\n *\n" );
34
35     for( i = 0; i < nrows ; i++ )
36     {
37         if( !accept[i].string )
38             fprintf(fp, " * State %d [nonaccepting]", i );
39         else
40         {
41             fprintf(fp, " * State %d [accepting, line %d <",
42                     i , ((int *) (accept[i].string)) [-1] );
43
44             fputstr( accept[i].string, 20, fp );
45             fprintf(fp, ">]" );
46

```



Listing 2.51. continued...

```

47         if( accept[i].anchor )
48             fprintf( fp, " Anchor: %s%s",
49                         accept[i].anchor & START ? "start" : "",
50                         accept[i].anchor & END   ? "end"   : "" );
51     }
52
53     last_transition = -1;
54     for( j = 0; j < MAX_CHARS; j++ )
55     {
56         if( dtran[i][j] != F )
57         {
58             if( dtran[i][j] != last_transition )
59             {
60                 fprintf(fp, "\n *      goto %2d on ", dtran[i][j] );
61                 chars_printed = 0;
62             }
63
64             fprintf(fp, "%s", bin_to_ascii(j,1) );
65
66             if( (chars_printed += strlen(bin_to_ascii(j,1))) > 56 )
67             {
68                 fprintf(fp, "\n *          " );
69                 chars_printed = 0;
70             }
71
72             last_transition = dtran[i][j];
73         }
74     }
75     fprintf(fp, "\n");
76 }
77 fprintf(fp," */\n\n" );
78 fprintf(fp,"#endif\n" );
79 }
80
81 /*-----*/
82
83 PUBLIC void pdriver( output, nrows, accept )
84 FILE    *output;
85 int      nrows;           /* Number of states in dtran[]           */
86 ACCEPT   *accept;         /* Set of accept states in dtran[]       */
87 {
88     /* Print the array of accepting states, the driver itself, and the case
89     * statements for the accepting strings.
90     */
91
92     int      i;
93     static char *text[] =
94     {
95         "The Yyaccept array has two purposes. If Yyaccept[i] is 0 then state",
96         "i is nonaccepting. If it's nonzero then the number determines whether",
97         "the string is anchored, 1=anchored at start of line, 2=at end of",
98         "line, 3=both, 4=line not anchored",
99         NULL
100    };
101
102    comment( output, text );
103    fprintf(output, "YYPRIVATE YY_TTYPE Yyaccept[] =\n" );
104    fprintf(output, "{\n" );
105

```



Listing 2.51. continued...

```

106     for( i = 0 ; i < nrows ; i++ )                                /* accepting array */
107     {
108         if( !accept[i].string )
109             fprintf( output, "\t0  " );
110         else
111             fprintf( output, "\t%-3d", accept[i].anchor ? accept[i].anchor :4);
112
113         fprintf(output, "%c      /* State %-3d */\n",
114                         i == (nrows -1) ? ' ' : ',', i );
115     }
116     fprintf(output, "};\n\n" );
117
118     driver_2( output, !No_lines );                                     /* code above cases */
119
120     for( i = 0 ; i < nrows ; i++ )                                     /* case statements */
121     {
122         if( accept[i].string )
123         {
124             fprintf(output, "\t\tcase %d:\t\t\t/* State %-3d */\n",i,i);
125             if( !No_lines )
126                 fprintf(output, "#line %d \"%s\"\n",
127                               *( int *) (accept[i].string) - 1,
128                               Input_file_name );
129
130             fprintf(output, "\t\t\t%s\n",    accept[i].string   );
131             fprintf(output, "\t\t\tbreak;\n"           );
132         }
133     }
134
135     driver_2( output, !No_lines );                                     /* code below cases */
136 }
```

Listing 2.52. *lex.c*—main() and Other High-Level Functions (Part 1)

```

1  #include <stdio.h>
2  #include <ctype.h>
3  #include <stdarg.h>
4  #include <tools/debug.h>                                         /* For vprintf() */
5  #include <tools/set.h>
6  #include <tools/hash.h>
7  #include <tools/compiler.h>
8  #include "dfa.h"
9
10 #define ALLOCATE
11 #include "globals.h"
12
13 /*-----*/
14
15 PUBLIC void error      (int usage, char *fmt, ...);
16 PUBLIC int main        (int argc, char **argv);
17 PRIVATE int getline     (char **stringp, int n, FILE *stream);
18 PRIVATE void do_file    (void);
19 PRIVATE void tail       (void);
20 extern char *get_expr  (void);                                     /* in input.c */
21
```



Listing 2.52. continued...

```

22 /*-----*/
23
24 #define DTRAN_NAME "Yy_nxt" /* Name used for DFA transition table. Up to */
25 /* 3 characters are appended to the end of */
26 /* this name in the row-compressed tables. */
27
28 #define E(x) fprintf(stderr,"%s\n", x)
29
30 PRIVATE int Column_compress = 1; /* Variables for command-line switches */
31 PRIVATE int No_compression = 0;
32 PRIVATE int Threshold = 4;
33 PRIVATE int No_header = 0;
34 PRIVATE int Header_only = 0;
35
36 extern int Verbose; /* in globals.h */
37 extern int No_lines; /* in globals.h */
38 extern int Lineno; /* In globals.h, the line number */
39 /* used to print #line directives. */
40
41 /*-----*/
42
43 PRIVATE void cmd_line_error( usage, fmt, ... )
44 char *fmt;
45 {
46     /* Print an error message and exit to the operating system. This routine is
47     * used much like printf(), except that it has an extra first argument.
48     * If "usage" is 0, an error message associated with the current value of
49     * errno is printed after printing the format/arg string in the normal way.
50     * If "usage" is nonzero, a list of legal command-line switches is printed.
51     *
52     * I'm using the ANSI conventions for a subroutine with a variable number of
53     * arguments. These differ a little from the Unix V conventions. Look up
54     * _doprnt() [it's in the "printf(3S)" entry] in the Unix Programmer's
55     * Manual for details. A version of the <stdarg.h> macros used here is in
56     * Appendix A.
57 */
58
59 extern char *sys_errlist[];
60 extern int errno;
61 va_list args;
62
63 va_start(args, fmt);
64 fprintf ( stderr, "lex: " );
65 vfprintf( stderr, fmt, args );
66
67 if( !usage )
68     fprintf( stderr, "(%s)\n", sys_errlist[errno] );
69 else
70 {
71     E("\n\nUsage is:    LeX [options] file" );
72     E("-f for (f)ast. Don't compress tables." );
73     E("-h suppress (h)eader comment that describes state machine." );
74     E("-H print the (H)eader only." );
75     E("-l Suppress #(l)ine directives in the output." );
76     E("-mS Use string S as template name rather than lex.par." );
77     E("-cN use pair (c)oression, N = threshold (default 4)." );
78     E("-t Send output to standard output instead of lex.yy.c" );
79     E("-u UNIX mode, newline is \\n, not \\n or \\r" );
80     E("-v (v)erbose mode, print statistics." );
81     E("-V More (V)erbose, print internal diagnostics as lex runs." );

```



Listing 2.52. continued...

```
82      }
83      exit(1);
84      va_end(args);
85  }
86
87 /*-----*/
88
89 PUBLIC void lerror( status, fmt, ... )
90 char *fmt;
91 {
92     /* Print an error message and input line number. Exit with
93      * indicated status if "status" is nonzero.
94      */
95
96     extern int errno;
97     va_list args;
98
99     va_start(args, fmt);
100    fprintf ( stderr, "lex, input line %d: ", Actual_lineno );
101    vfprintf ( stderr, fmt, args );
102    if( status )
103        exit( status );
104    va_end(args);
105 }
106
107 /*-----*/
108
109 PUBLIC main( argc, argv )
110 char **argv;
111 {
112     static char      *p ;
113     static int   use_stdout = 0;
114
115     signon();
116
117     for( ++argv, --argc; argc && *(p = *argv) == '-'; ++argv, --argc )
118     {
119         while( *++p )
120         {
121             switch( *p )
122             {
123                 case 'f': No_compression = 1; break;
124                 case 'h': No_header      = 1; break;
125                 case 'H': Header_only    = 1; break;
126                 case 'l': No_lines       = 1; break;
127                 case 'm': Template       = p + 1; goto out;
128                 case 'p': Public         = 1; break;
129                 case 't': use_stdout     = 1; break;
130                 case 'u': Unix           = 1; break;
131                 case 'v': Verbose        = 1; break;
132                 case 'V': Verbose        = 2; break;
133                 case 'c': Column_compress = 0;
134
135                 if( !isdigit(p[1]) )
136                     Threshold = 4;
137                 else
138                 {
139                     Threshold = atoi( ++p );
```



Listing 2.52. continued...

```

140                                while( *p && isdigit( p[1] ) )
141                                    ++p;
142                                }
143                                break;
144
145        default:    cmd_line_error(1, "-%c illegal argument.", *p);
146        break;
147    }
148}
149out:;
150}
151
152if( argc > 1 )
153    cmd_line_error( 1, "Too many arguments. Only one file name permitted" );
154
155else if( argc <= 0 )
156    cmd_line_error( 1, "File name required" );
157else /* argc == 1 */
158{
159    if( Ifile = fopen(*argv,"r") )
160        Input_file_name = *argv;
161    else
162        cmd_line_error( 0, "Can't open input file %s", *argv );
163}
164
165if( !use_stdout )
166    if( !(Ofile = fopen( Header_only ? "lexyy.h" : "lexyy.c", "w")) )
167        cmd_line_error( 0, "Can't open output file lexyy.[ch]" );
168
169do_file();
170fclose( Ofile );
171fclose( Ifile );
172exit( 0 );
173}
174
175/*-----*/
176
177PRIVATE void do_file( )
178{
179    int      nstates;           /* Number of DFA states          */
180    ROW     *dtran;            /* Transition table             */
181    ACCEPT  *accept;           /* Set of accept states in dfa */
182    FILE    *input, *driver_1(); /* Template file for driver    */
183    int      i;
184
185    /* Process the input file */
186
187    head( Header_only );       /* print everything up to first %% */
188
189    nstates = min_dfa( get_expr, &dtran, &accept );           /* make DFA */
190    if( Verbose )
191    {
192        printf("%d out of %d DFA states in minimized machine\n", nstates,
193                           DFA_MAX );
194        printf("%d bytes required for minimized tables\n\n",
195               nstates * MAX_CHARS * sizeof(TTYPE)                /* dtran */
196               + nstates * sizeof(TTYPE) );                      /* accept */
197    }
198}

```



Listing 2.52. continued...

```

199     if( !No_header )
200         pheader( Ofile, dtran, nstates, accept );           /* print header */
201                                         /* comment. */ */
202     if( !Header_only )
203     {
204         /* first part */
205         /* of driver. */
206     if( !(input = driver_1( Ofile, !No_lines, Template )) )
207     {
208         perror( Template );
209         exit(1);
210     }
211
212     if( No_compression )                                     /* compressed tables */
213     {
214         fprintf (Ofile , "YYPRIVATE YY_TTYPE %s[ %d ][ %d ] =\n",
215                               DTRAN_NAME, nstates, MAX_CHARS);
216
217         print_array( Ofile, (int *)dtran, nstates, MAX_CHARS );
218         defnext    ( Ofile, DTRAN_NAME );
219     }
220     else if( Column_compress )                                /* column-compressed tables. */
221     {
222         i = squash ( Ofile, dtran, nstates, MAX_CHARS, DTRAN_NAME );
223         cnext      ( Ofile, DTRAN_NAME );
224
225         if( Verbose )
226             printf("%d bytes required for column-compressed tables\n\n",
227                     i                                         /* dtran */
228                     + (nstates * sizeof(int)) );           /* Yy_accept */
229     }
230     else                                                       /* pair-compressed tables */
231     {
232         i = pairs( Ofile, (int *)dtran, nstates,
233                         MAX_CHARS, DTRAN_NAME, Threshold, 0 );
234         if( Verbose )
235         {
236             /* Figure the space occupied for the various tables. The
237             * Microsoft compiler uses roughly 100 bytes for the yy_next()
238             * subroutine. Column compression does yy_next in line with a
239             * macro so the overhead is negligible.
240             */
241
242             i =   (i          * sizeof(TTYPE))    /* YysDD arrays      */
243                 + (nstates * sizeof(TTYPE*))   /* Dtran[]           */
244                 + (nstates * sizeof(TTYPE))   /* Yy_accept[]      */
245                 + 100                      ; /* yy_next()        */
246
247             printf("%d bytes required for pair-compressed tables\n", i );
248         }
249
250         pnext( Ofile, DTRAN_NAME );
251     }
252
253     pdriver( Ofile, nstates, accept );           /* print rest of driver */
254                                         /* and everything following */
255                                         /* the second %% */
256 }
257 }
258 }
```



Listing 2.52. continued...

```

259  /* -----
260   * Head processes everything up to the first %%. Any lines that begin
261   * with white space or are surrounded by %{ and %} are passed to the
262   * output. All other lines are assumed to be macro definitions.
263   * A %% can not be concealed in a %{ %} but it must be anchored at start
264   * of line so a %% in a printf statement (for example) is passed to the
265   * output correctly. Similarly, a %{ and %} must be the first two characters
266   * on the line.
267 */
268
269 PRIVATE head( suppress_output )
270 int suppress_output;
271 {
272     int transparent = 0;           /* True if in a %{ %} block */
273
274     if( !suppress_output && Public )
275         fputs( "#define YYPRIVATE\n\n", Ofile );
276
277     if( !No_lines )
278         fprintf( Ofile, "#line 1 \"%s\"\n", Input_file_name );
279
280     while( fgets( Input_buf, MAXINP, Ifile ) )
281     {
282         ++ Actual_lineno;
283
284         if( !transparent )           /* Don't strip comments */
285             strip_comments( Input_buf ); /* from code blocks. */
286
287         if( Verbose > 1 )
288             printf( "h%d: %s", Actual_lineno, Input_buf );
289
290         if( Input_buf[0] == '%' )
291         {
292             if( Input_buf[1] == '%' )
293             {
294                 if( !suppress_output )
295                     fputs( "\n", Ofile );
296                 break;
297             }
298             else
299             {
300                 if( Input_buf[1] == '{' )           /* }{*/
301                     transparent = 1;
302
303                 else if( Input_buf[1] == '}' )
304                     transparent = 0;
305
306                 else
307                     lerror(0, "Ignoring illegal %%c directive\n",
308                             Input_buf[1] );
309             }
310         }
311         else if( transparent || isspace( Input_buf[0] ) )
312         {
313             if( !suppress_output )
314                 fputs( Input_buf, Ofile );
315         }
316         else
317         {
318             new_macro( Input_buf );

```



Listing 2.52. continued...

```
319         if( !suppress_output )
320             fputs( "\n", Ofile );    /* Replace macro def with a blank */
321                                         /* line so that the line numbers */
322                                         /* won't get messed up.          */
323     }
324 }
325
326 if( Verbose > 1 )
327     printmacs();
328 }
329
330 /*-----*/
331
332 strip_comments( string )
333 char *string;
334 {
335     /* Scan through the string, replacing C-like comments with space
336      * characters. Multiple-line comments are supported.
337      */
338
339     static int incomment = 0;
340
341     for( *string ; ++string )
342     {
343         if( incomment )
344         {
345             if( string[0]=='*' & string[1]=='/' )
346             {
347                 incomment = 0 ;
348                 *string++ = ' ' ;
349             }
350
351             if( !isspace( *string ) )
352                 *string = ' ' ;
353         }
354         else
355         {
356             if( string[0]=='/' & string[1]=='*' )
357             {
358                 incomment = 1 ;
359                 *string++ = ' ' ;
360                 *string = ' ' ;
361             }
362         }
363     }
364 }
365
366 /*-----*/
367
368 PRIVATE void tail()
369 {
370     fgets(Input_buf, MAXINP, Ifile);    /* Throw away the line that      */
371                                         /* had the %% on it.            */
372
373     if( !No_lines )
374         fprintf( Ofile, "#line %d \"%s\"\n",
375                  Actual_lineno + 1, Input_file_name);
376 }
```



Listing 2.52. continued...

```

377     while( fgets(Input_buf, MAXINP, Ifile) )
378     {
379         if( Verbose > 1 )
380             printf( "t%d: %s", Actual_lineno++, Input_buf );
381
382         fputs(Input_buf, Ofile);
383     }
384 }
```

2.6 Exercises

- 2.1. Modify the simple expression grammar used in Chapter 1 so that COMMENT tokens can appear anywhere in the expression. This exercise demonstrates graphically why it's useful for a lexical analyzer to do things like remove comments from the input.
- 2.2. Write a lexical analyzer that processes the PL/I statement

if then then then = else; else else = then;
correctly, distinguishing the keywords from the identifiers.

- 2.3. Write a regular expression that can recognize all possible integer constants in C. All of the following should be recognized:

0x89ab 0123 45 'z' '\t' '\xab' '\012'

- 2.4. Write a regular expression that recognizes a C identifier. Construct an NFA for this expression using Thompson's Construction, convert the NFA to a DFA using Subset Construction, and minimize the DFA.
- 2.5. Write a set of input functions that can handle interactive input from a human being more elegantly than the ones in *input.c*. (They should support command-line editing, and so forth). Your routines must work with the default driver generated by L^EX (they must use the same interface as the default ones). The routines should be able to get input directly from the console, rather than from standard input.
- 2.6. Use L^EX to rewrite its own lexical analyzer.
- 2.7. Write a set of regular expressions that can recognize C subroutine declarations. You can assume that the entire declaration is on one line. Declarations such as the following should be recognized:

```
TYPE          *doo( wha, ditty )
unsigned char (*foo()) [10]
int           ( *signal(signal, ptr) )()
```

but not the following

```
extern int   ( *signal(int signal, (*ptr)()) )();
sam( dave );
```

- 2.8. Use the regular expressions developed in the previous exercise to write a L^EX program that inputs a list of C programs whose names are listed on the command line, and which outputs all subroutine declarations along with the name of the file in which they are found and their input line numbers.

- 2.9. Modify the previous exercise so that only the subroutine name is printed, rather than the entire declaration.
- 2.10. Various UNIX editors support the \ (and \) operators, which bracket subexpressions in a regular expression. For example, an expression like this:

```
\(abc\)\def\(ghi\)
```

recognizes the string abcdefghi, and the two substrings abc and ghi can be accessed later using the notation \1 and \2 respectively. Modify LEX to support a similar feature, but as follows:

A subexpression in a regular expression is indicated with parenthesis.

- a. Modify LEX so that repeated subexpressions in the regular expressions themselves can be referenced using the notation \0, \1, and so forth. \0 references the first subexpression, \1 the next one, and so forth. A subexpression's number is derived by counting the number of open parentheses to the left of the one that starts that subexpression. For example, the following regular expression recognizes three, comma-delimited decimal numbers:

```
[0-9]+, [0-9]+, [0-9]+
```

You should be able to express the same thing as follows:

```
([0-9]+) (,) \0\1\0
```

The \0 in the regular expression tells LEX to duplicate the first subexpression ([0-9]+), the \1 duplicates the second (,).

The double-quote character and double backslash should still work as expected. Both of the following expressions should recognize an input string like "1234\1\2\1":

```
([0-9]+) (,) "\1\2\1"
([0-9]+) (,) \\1\\2\\1
```

The backslashes are treated as literal characters here.

- b. Modify LEX so that the lexemes associated with the matched subexpression can be accessed in an action using an argv-like array of string pointers called yyv: yyv[0] holds the lexeme that matches the first subexpression, yyv[1] the second, and so on. An integer called yyc should hold the number of valid entries in yyv. For example, the expression in the following rule recognizes the input pattern abcdefghi:

```
((abc(def)) (ghi))      while( --yyc >= 0 )
                           printf("%s\n", *yyv++ );
```

The associated action prints:

```
abcdefghi
abcdef
def
ghi
```

- 2.11. Modify LEX to support a unary NOT operator. It should be higher precedence than the closure operators and lower precedence than concatenation. It should precede its operand. For example, !c matches anything except the character c; the expression a! (bc) should match an a followed by any sequence of characters other than the two character sequence bc; !(and|or) should match any string except and or or; finally, ! [a-z] should work just like [^a-z]. Be careful of the newline, which should not be matched by any of the foregoing patterns.

- 2.12. Add the `/` lookahead operator to L^EX. The expression a/b — where a and b are both regular expressions — matches a , but only if it is followed by b . The characters that match b should not be put into the lexeme. In fact, they should be pushed back into the input after being recognized. For example, the regular expression abc/def should match an abc only if it's followed by a def , and the def should be pushed back into the input for subsequent processing. You can simulate the `$` operator with `xxx/[\r\n]`. Similarly, the following pattern matches the string *if*, but only if the next nonwhite character isn't a semicolon, plus sign, minus sign, or asterisk:

```
if/[\s\t\n\r]*[^;+-*]`
```

By the same token, `if/[\s\t\n\r]*(`` matches an *if* only if the next nonwhite character is an open parenthesis.

- 2.13. Add the REJECT directive to L^EX. This directive selects the next alternative accepting action when regular expressions conflict. For example, the following recognizes the word *he*, even if it's imbedded in *she*:

```
int he=0, she=0;
%%
he      { ++he;  REJECT; }
she     { ++she; REJECT; }

(.|\n) ; /* ignore */
%%
main()
{
    yylex();
    printf("Saw %d he's and %d she's\n", he, she );
}
```

It recognizes the *she* first, then the REJECT action causes it to execute the code associated with the *he* regular expressions too. Note that it's not doing anything special with the input, it's only executing all possible accepting actions for the expression. The program should print

Saw 2 he's and 1 she's

when given the input:

he she

- 2.14. Add the $X\{a,b\}$ operator to L^EX. This expression matches X repeated a to b times. For example `l{1,3}ama` recognizes *lama*, *llama*, and *llllama*.
- 2.15. Add left-context sensitivity to L^EX. Do it like this: The statement

```
%START NAME
```

is placed in the first, definitions part of the L^EX input file. It creates a class of regular expressions that are active only in certain situations (discussed below). NAME can be an arbitrary name having up to eight characters. All regular expressions that start with `<NAME>` are placed in the class. The angle brackets must be present, like this:

```
<NAME>(a|b)* { some_code(); }
```

Activate all states in the class with the following line anywhere in the code part of a rule:

```
BEGIN NAME;
```

That is, the earlier regular expression, $(a/b)^*$, won't be recognized until a BEGIN

NAME has been executed in a code block. Use BEGIN 0; to disable all special regular expressions. An expression can be put into several classes like this:

```
<NAME1,NAME2,NAME3>expression { some_code(); }
```

Once an expression is activated, the usual input-file precedence rules are also active. That is, if a special expression occurs earlier in the input file than a conflicting expression, the special one takes precedence once it's activated.

- 2.16. Modify L^EX to detect which of the two compression methods yields the smallest tables, and then use that method automatically.
- 2.17. Add a command-line switch to L^EX that causes the two compression methods to be combined: Pair compression should be applied first, then a map array should be used to compress the uncompressed rows in the pair-compressed tables. Only one copy of two or more identical rows should be kept, with several pointers in Yy_nxt [] addressing the single copy of the row.
- 2.18. Modify the minimization routines to remove dead states—nonaccepting states that have no outgoing transitions other than ones that loop back into the current state.
- 2.19. Modify the minimization routines to remove unreachable states—states that have no incoming edges.
- 2.20. Sometimes it's useful for a program's users to be able to modify the token set when they install a program. This way, they can customize a program's input syntax to fit their own preferences. Modify L^EX so that it produces a file containing binary representations of the state and auxiliary tables rather than C source code. Modify yylex() so that it reads the tables from this file into memory the first time it is called. This way, the user can change the L^EX input file and generate a new state table without having to recompile the program that is using that state table.

Note that the actions still have to be put into the compiled program—there's no way to read them from disk at run time. Consequently, though the regular expressions can be modified at installation time, the number of regular expressions must stay the same. The position of the regular expression from the top of the file determines the action to execute when that expression is recognized. That is, the second action in the original input file is executed when the second regular expression is matched, even if that expression is changed at installation time.

Context-Free Grammars

The discussion of formal grammars started in Chapter One continues here. I discuss the general, more theoretical aspects of grammars, but with emphasis on building real grammars, discussing precedence and associativity, list and expression grammars, and so forth. Attributed and augmented grammars are introduced as well, and they are put into the context of a recursive-descent parser. Grammars are discussed further in Chapter Four, where various grammatical transformations are presented.

3.1 Sentences, Phrases, and Context-Free Grammars

This section quickly summarizes the terms introduced in Chapters One and Two and extends these terms somewhat.

An *alphabet* is a collection of characters. The ASCII character set is a good example of an alphabet. A *string* or *word* is a specific sequence of symbols from an input *alphabet*. In compiler parlance, a word in the language is called a *token*—an indivisible lexical unit of the language.

A *language* is a set of words, and a *sentence* is a sequence of one or more of the words that comprise the language. Strictly speaking, there is no inherent ordering of the words within the sentence, and that's where a *grammar* comes into play. A formal *grammar* is a system of rules (called *productions*) that control the order in which words may occur in a sentence.

The *syntax* of a sentence determines the relationships between the words and phrases in a sentence. That is, the syntax of a language controls the structure of a sentence and nothing more. Taking English as an example, the syntax of English requires that the subject precede the verb, or that an adjective must precede the noun that it modifies. The syntax rules tell you nothing about what the sentence actually means, however.

Informally, a *context-free grammar* is a system of definitions that can be used to break up a sentence into phrases solely on the basis of the sequence of strings in the input sentence. A context-free grammar is usually represented in *Backus-Naur form* (BNF), in which productions are represented as follows:

part_of_speech → *definition*

The → operator (often represented as ::=) is the “is defined as” or “goes to” operator.

Alphabet.

String, word.

Token.

Language, sentence.

Grammar, production.

Syntax.

Context-free grammar.

Backus-Naur Form, →,
::=.

Left-hand side (LHS).

The left-hand side of the production (often abbreviated to LHS), is the name of the syntactic unit within the language (noun, verb, prepositional phrase, and so forth). The right-hand side (or RHS) is a list of symbols that define the LHS. For example:

Right-hand side (RHS).

sentence → *subject predicate PERIOD*

A *sentence* is a *subject*, followed by a *predicate*, followed by a **PERIOD**. There are two types of symbols on the right-hand side: *Nonterminal* symbols are terms that have to be defined further somewhere in the grammar (they have to appear on a left-hand side), and *terminal* symbols such as tokens which need no further definition. The nonterminal nodes form interior nodes in the parse tree; the terminal nodes are all at the ends of branches. There are terminal symbols other than tokens, discussed below.

So, a context-free grammar is one in which all input sentences can be parsed strictly on the basis of syntax. Formally, a context-free grammar is composed of the following:

Nonterminal and terminal symbols.

Context-free grammar: formal definition.

- A finite set of terminal symbols or tokens.
- A finite set of nonterminal symbols.
- A finite set of productions of the form $s \rightarrow \alpha$, where s is a nonterminal symbol, and α is a list of zero or more terminal and nonterminal symbols. s is called the *left-hand side* of the production and α is called the *right-hand side*. Every nonterminal symbol that appears in the right-hand side in some production must also appear on a left-hand side. No terminal symbols may appear on a left-hand side.
- A single *start* or *goal* symbol from which all the productions derive.

It's theoretically possible to have a “context-sensitive” grammar that does consider semantics when analyzing a sentence, but this type of grammar is beyond the scope of this book and won't be discussed further.

A grammar can also distinguish groups of words that are related syntactically. For example, in English, the syntax of the language helps separate a subject from a predicate, recognize prepositional phrases, and so forth. Each of these syntactically related lists of words are called *phrases*. In a computer language, a loop-control statement and associated body is a phrase, as is an entire subroutine declaration. Note that the phrases can be organized hierarchically—a sentence can contain a predicate, which can in turn contain a prepositional phrase. A grammatical phrase can be as short as a single word: in C, a single identifier can comprise a phrase. It can also be empty—I'll discuss this case below. A production is a rule that describes the syntax of a single phrase.

Phrase.

Syntactic rules are used to break a sentence into its component parts of speech and analyze the relationship of one part to another (this process is called *parsing* the sentence). The term is used in both linguistic and compiler theory. A *parser* is a computer program that uses a context-free grammar to parse an input sentence—to isolate the component parts of speech for subsequent processing. Most parser programs are also *recognizer* programs. They accept (returns *yes*) only if the input forms a syntactically correct sentence. The fact that a parser can also generate code is actually immaterial from a theoretical point of view.

Parsing.

The parser can analyze only the structure of the sentence. Even in English, the syntax of a sentence tells you nothing about what the sentence means. Meaning comes under the purview of *semantics*. A context-free grammar defines the structure of the sentence only. It tells you nothing about the semantics—the meanings of the words—except that you can sometimes infer meaning from the position of the word within a sentence. For example, a grammar could tell you that a variable declaration is comprised of a sequence of type-specification tokens (**int**, **long**, **short**, and so forth) followed by a single identifier, but it is just manipulating sequences of strings, it has no idea what the strings mean. In the case of a declaration, it knows nothing about the characteristics of the individual types—as far as the grammar is concerned, the following declarations are

Recognizer programs.

Semantics and syntax.

all reasonable because they're all sequences of type specifiers followed by a single identifier:

```
long      john_silver;
short     stop;
long short muddle;
```

What the grammar *does* do is distinguish the specifiers from the identifiers. If a type specifier were redefined so that it could represent an **int** only, then the meaning of the specifier could be extrapolated from the syntax. That is, when the parser encountered a type specifier, it would know that the specifier was an **int**, because that was the only possibility. This sort of degenerate case is not much use in natural languages like English, but it can be used effectively in languages with limited syntactic and semantic scope, such as programming languages, which might have only a few verbs (**while**, **do**, **for**), and the same number of nouns (**int**, **float**, **double**).

Syntax-directed translation.

In *syntax-directed translation*, code is generated (the input language is translated to the output language) strictly on the basis of the syntax of a sentence—the meaning of a phrase or word is inferred from its position within the sentence. In practical terms, the parser performs certain code-generation actions as it traverses a parse tree which is created using syntactic rules. It performs a code-generation action every time a node on the tree is visited.

3.2 Derivations and Sentential Forms

Derivations.

A *derivation* is a way of showing how an input sentence can be recognized with a grammar. You can also look at it as a type of algebraic substitution using the grammar. As an example of the derivation process, you start a leftmost derivation with the topmost nonterminal in a grammar, called the *start* or *goal* symbol. You then replace this nonterminal with one of its right-hand sides. Continue in this manner—but replace the leftmost nonterminal created from the previous step with one of the right-hand sides of that nonterminal—until there are no more nonterminals to replace.

The following grammar recognizes a simple semicolon-terminated expression consisting of interspersed numbers and **PLUS** tokens:

1:	<i>stmt</i>	\rightarrow	<i>expr SEMI</i>
2:	<i>expr</i>	\rightarrow	<i>factor PLUS expr</i>
3:			<i>factor</i>
4:	<i>factor</i>	\rightarrow	NUMBER

The \Rightarrow symbol.

All you need do to prove that an input sentence (like `1+2;1+2;`) can be recognized for this grammar is come up with a derivation that can match the input sentence, token for token. The following derivation does just that (the \Rightarrow symbol means *derives*),

<i>stmt</i>	\Rightarrow	<i>expr SEMI</i>	<i>by Production 1</i>
	\Rightarrow	<i>factor PLUS expr SEMI</i>	<i>by Production 2</i>
	\Rightarrow	<i>NUMBER PLUS expr SEMI</i>	<i>by Production 4</i>
	\Rightarrow	<i>NUMBER PLUS factor SEMI</i>	<i>by Production 3</i>
	\Rightarrow	<i>NUMBER PLUS NUMBER SEMI</i>	<i>by Production 4</i>

Leftmost derivation.

This process is called a *leftmost* derivation because you always replace the leftmost nonterminal in the partially-parsed sentence with the equivalent production's right-hand side. Start with the goal symbol (the topmost symbol in the grammar) and then replace the leftmost nonterminal at each step. The \xrightarrow{L} operator signifies a leftmost derivation. $X \xrightarrow{L} Y$ means “X derives Y by a leftmost derivation.” You can say that

$\text{expr SEMI} \stackrel{L}{\Rightarrow} \text{factor PLUS expr SEMI}$

by substituting *expr* for its right-hand side. The foregoing derivation was created in a *top-down* fashion, starting with the topmost symbol in the grammar and working down to the tokens. This process is called a *top-down parse*.

Each of the steps in the earlier example (each partial derivation) is formally called a *viable prefix*. You can look at a viable prefix as a partially parsed sentence. Another useful term is a *handle*, which is the portion of the viable prefix that is replaced at each step in the derivation. The handle is always a contiguous sequence of zero or more symbols at the far right or far left of the viable prefix (depending on how the derivation is done).

You can also have a *rightmost derivation* which starts with the goal symbol and replaces the rightmost nonterminal with each step. The $\stackrel{R}{\Rightarrow}$ operator can be used to indicate this kind of derivation. It's also possible to have a *bottom-up* parse—one that starts with the leaves of the tree and works its way up to the root. Bottom-up parsers use right-most derivations. They start with the input symbols and, when a *right-hand side* is encountered in a viable prefix, replaces the symbols that comprise the right-hand side with the equivalent left-hand side. You still go through the input from left to right, but you always replace the rightmost handle in the viable prefix. Using the same grammar and input as before, you can derive the goal symbol from the input symbols as follows:

$\text{NUMBER PLUS NUMBER SEMI}$ $\stackrel{R}{\Rightarrow} \text{factor PLUS NUMBER SEMI}$ $\stackrel{R}{\Rightarrow} \text{factor PLUS factor SEMI}$ $\stackrel{R}{\Rightarrow} \text{factor PLUS } \stackrel{R}{\Rightarrow} \text{expr SEMI}$ $\stackrel{R}{\Rightarrow} \text{expr SEMI}$ $\stackrel{R}{\Rightarrow} \text{stmt}$	<i>Start with the input.</i> <i>Apply factor</i> \rightarrow NUMBER <i>Apply factor</i> \rightarrow NUMBER <i>Apply expr</i> \rightarrow <i>factor</i> <i>Apply expr</i> \rightarrow <i>factor</i> PLUS <i>expr</i> <i>Apply stmt</i> \rightarrow <i>expr SEMI</i>
---	--

The parser scans the viable prefix (the partially-parsed sentence) for a handle that matches the right-hand side of some production, and then replaces the handle with the left-hand side of that production. For example, **NUMBER** is the right-hand side of *factor* \rightarrow **NUMBER**, and this production is applied in the first derivation, replacing the **NUMBER** with the *factor*. *factor PLUS* doesn't form a right hand side, so the parser moves to the next symbol, the second **NUMBER**, for the next replacement, and so forth. A practical bottom-up parser, which uses a rightmost derivation, is discussed in great detail in Chapter Five, so don't worry about it if the foregoing didn't sink in.

Note that top-down parsers that do rightmost derivations and bottom-up parsers that do leftmost derivations are theoretically possible, but they are difficult to implement. Similarly, you can go through the input from right to left rather than left to right, but again, the practice is not common.

In general, the \Rightarrow , $\stackrel{L}{\Rightarrow}$, and $\stackrel{R}{\Rightarrow}$ symbols mean that only one substitution has occurred between the left- and right-hand side of the equation. Several other forms of the \Rightarrow symbol are commonly used: $\stackrel{*}{\Rightarrow}$ and $\stackrel{+}{\Rightarrow}$ means derives in zero or more and in one or more steps, respectively. The various forms can also be combined in reasonable ways: $\stackrel{L}{\Rightarrow} \stackrel{R}{\Rightarrow}$, $\stackrel{*}{\Rightarrow} \stackrel{+}{\Rightarrow}$.

We need one more term before continuing: if $S \stackrel{L}{\Rightarrow}_* \alpha$ (if there is a leftmost derivation of the viable prefix α from the nonterminal S), then α is sometimes called a *left-sentential form* of S . *Right-sentential form* is used for a rightmost derivation and a just plain *sentential form* is used when the type of derivation isn't known or doesn't matter. For example, **NUMBER PLUS expr SEMI** is a left-sentential form of *stmt* because $\text{stmt} \stackrel{L}{\Rightarrow}_* \text{NUMBER PLUS expr SEMI}$.

Top-down parse.

Viable prefix, handle.

Rightmost derivation.
Bottom-up parse.

Sentential forms.

LL parsers and grammars.

LR parsers and grammars.

3.2.1 LL and LR Grammars

The concepts of left and right derivations apply directly to the two types of parsers that are discussed in Chapters Four and Five. In fact, the derivation closely follows the process that a parser uses to analyze the input sentence. Chapter Four discusses LL parsers, which go through the input stream from left to right (that's the first *L*), performing a leftmost derivation (that's the second *L*). Recursive-descent parsers, such as the ones examined in Chapter 1, are also LL parsers.

An LL grammar is a grammar that can be parsed by an LL parser. (There are many grammars that can't be so parsed.) An LL(1) grammar is a grammar that can be parsed by an LL parser with one symbol of lookahead. That is, if a nonterminal has more than one right-hand side, the parser can decide which one to apply by looking at the next input symbol.

It's possible to have an LL(*k*) grammar, where *k* is some number other than 1, and which requires *k* symbols of lookahead, but such grammars are not very practical. If the number is missing, 1 is assumed.

The other type of parser that we'll look at in this book is an LR parser which goes through the input from left to right, but does a bottom-up, rightmost derivation. The class of grammars that can be parsed by LR parsers are LR(*k*) grammars, and as before, we're interested primarily in LR(1) grammars, which require the parser to look ahead only one symbol to recognize a handle. LALR(*k*) grammars ("Look-Ahead" LR grammars), which are also discussed in depth in Chapter Five, are a special case of the more general LR(*k*) grammars.

3.3 Parse Trees and Semantic Difficulties

The grammar in Table 3.1 recognizes a very limited set of English sentences: A sentence is a subject and a predicate. Subjects are made up of nouns with an optional preceding article; predicates are either single verbs or a verb followed by a prepositional phrase; a prepositional phrase is a preposition followed by an object; and so on. A noun can be one of the tokens **time** or **arrow**.

You can use this grammar to parse the sentence: **Time flies like an arrow.** with the following leftmost derivation:

<i>sentence</i>	\xrightarrow{L}	<i>subject predicate</i>	<i>by Production 1</i>
	\xrightarrow{L}	<i>noun predicate</i>	<i>by Production 2</i>
	\xrightarrow{L}	time predicate	<i>by Production 9</i>
	\xrightarrow{L}	time verb prep_phrase	<i>by Production 5</i>
	\xrightarrow{L}	time flies prep_phrase	<i>by Production 13</i>
	\xrightarrow{L}	time flies preposition object	<i>by Production 6</i>
	\xrightarrow{L}	time flies like object	<i>by Production 14</i>
	\xrightarrow{L}	time flies like article noun	<i>by Production 8</i>
	\xrightarrow{L}	time flies like an noun	<i>by Production 12</i>
	\xrightarrow{L}	time flies like an arrow	<i>by Production 10</i>

The derivation process can be represented diagrammatically as a *parse tree*, a graph in which the root node of each subtree is a nonterminal symbol in the language, and the subtree itself represents the right-hand side of a replaced symbol. The goal symbol is always at the apex of the parse tree. The parse tree for the previous derivation is in Figure 3.1.

Note that the parser would happily have applied *noun*→**time** in the last step of the derivation, even though *time flies like a time* is not a reasonable sentence, and this is an example of a *semantic* versus a *syntactic* problem—a meaning-related versus a structural

Semantic difficulties in parsing.

Table 3.1. A Grammar for an English Sentence

Production Number	Production	
	<i>Left-hand Side</i>	<i>Right-hand Side</i>
1.	<i>sentence</i>	\rightarrow <i>subject predicate</i>
2.	<i>subject</i>	\rightarrow <i>noun</i>
3.		<i>article noun</i>
4.	<i>predicate</i>	\rightarrow <i>verb</i>
5.		<i>verb prep_phrase</i>
6.	<i>prep_phrase</i>	\rightarrow <i>preposition object</i>
7.	<i>object</i>	\rightarrow <i>noun</i>
8.		<i>article noun</i>
9.	<i>noun</i>	\rightarrow time
10.		arrow
11.	<i>article</i>	\rightarrow a
12.		an
13.	<i>verb</i>	\rightarrow flies
14.	<i>preposition</i>	\rightarrow like

problem. This situation often arises in computer-language implementations. The foregoing problem could probably be fixed grammatically, but it's not always possible to do so, and the resulting grammar will always be larger than necessary (the larger the grammar, the slower the parser). There's an even more serious problem that becomes clear when you add the following productions to the grammar:

$$\textit{noun} \rightarrow \textbf{fruit} \mid \textbf{banana}$$

The grammar now parses:

Fruit flies like a banana

without errors. Nonetheless, the sentence is actually parsed incorrectly because **flies** is being treated as a verb, not a noun. That is, you've just given bananas the capability of independent flight. You could try to rectify the situation by adding the following productions:

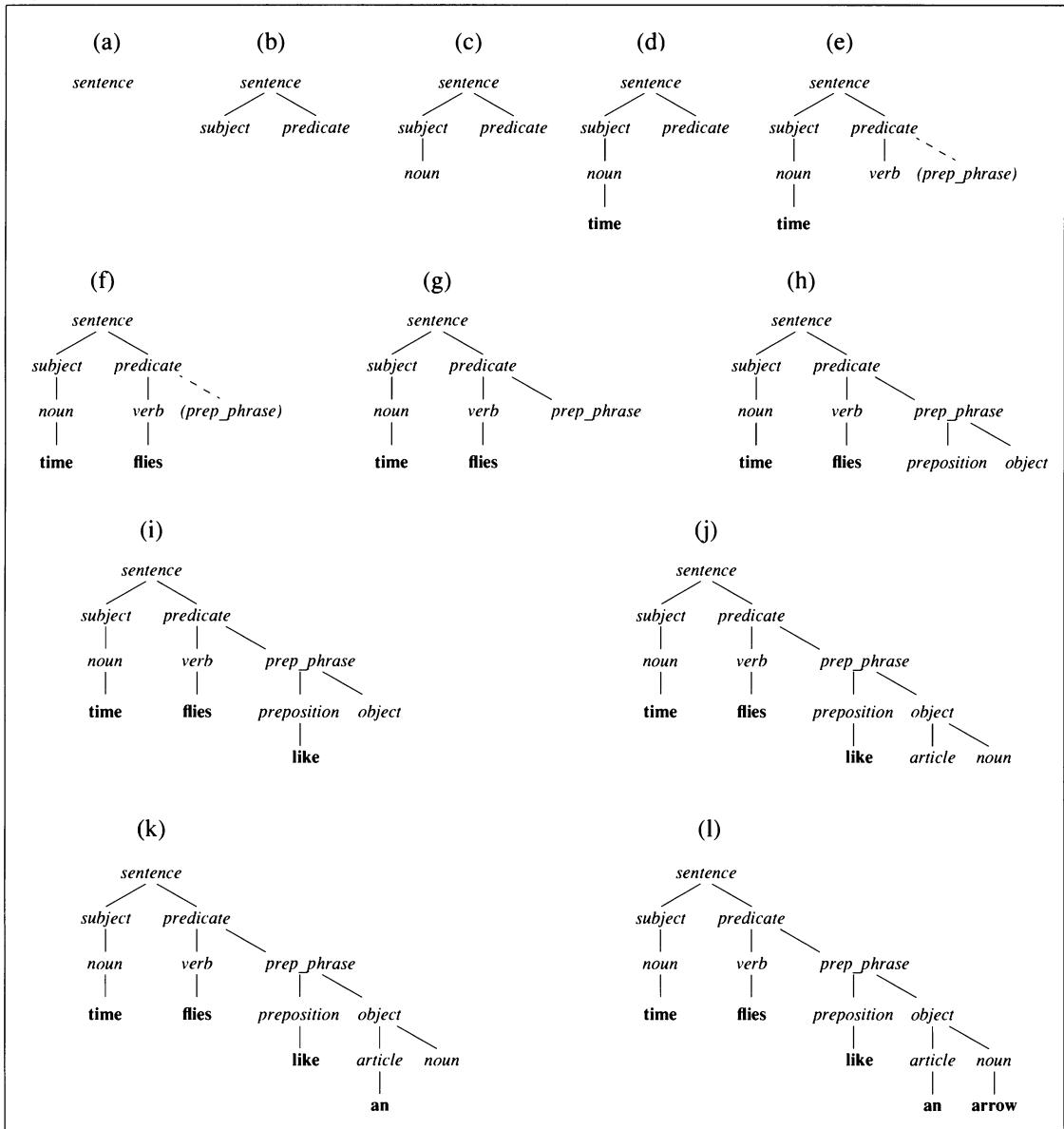
$$\begin{array}{lcl} \textit{adjective} & \rightarrow & \textbf{fruit} \\ \textit{noun} & \rightarrow & \textbf{flies} \end{array}$$

and modifying *subj* as follows:

$$\textit{subj} \rightarrow \textit{noun} \mid \textit{adjective noun}$$

But now, the grammar has become almost impossible to parse because the parser has to know the meanings of the various words in order to figure out which of the productions to apply. Making the wrong choice causes further errors (like the existence of a time fly, which can move forward or backwards in time to avoid being swatted). To parse both *time flies like an arrow* and *fruit flies like a banana* correctly, you'd have to know that there are no such things as time flies and that fruit can't fly, and it's very difficult (if not impossible) to build this semantic knowledge into a context-free grammar.

The foregoing example illustrates the limitations of context-free grammars, which can only specify syntax, not semantics. Though there are a few exceptions in artificial-intelligence applications, most real parsers are limited by the limitations of the grammar—they only know about syntax. Consequently, if semantic knowledge is

Figure 3.1. Evolution of a Parse Tree for *Time Flies Like An Arrow*

necessary to control the parse of a specific grammar, it becomes very difficult to build a computer program that can parse an input sentence. It's usually not possible to resolve all the semantic ambiguities of a computer language in the grammar. For example, the grammar itself can't tell you whether a variable is of the correct type when it's used in an expression, it can't recognize multiple declarations like `int x`, `long x;`, and so forth. To do the foregoing, you need to know something about meaning. Errors like the foregoing are best detected by the code-generation part of the compiler or by auxiliary code in the parser, not in the grammar itself.

3.4 ϵ Productions

I mentioned earlier that a grammatical phrase (a right-hand side) can be empty. More correctly, it's permitted for a right-hand side to match the empty string, ϵ . Productions of the form $s \rightarrow \epsilon$ are called ϵ *productions* (pronounced “epsilon productions”).

Consider the following grammar that recognizes compound statements:

1:	<i>compound_stmt</i>	\rightarrow	LEFT_CURLY <i>stmt</i> RIGHT_CURLY
2:	<i>stmt</i>	\rightarrow	NUMBER
3:			ϵ
	...		

I've simplified by defining a statement as either a single number or an empty string. This grammar supports empty compound statements, as does C. First, note that an ϵ production can effectively disappear from a derivation. The input {} generates the following derivation:

$$\begin{array}{lcl} \textit{compound_stmt} & \xrightarrow{\quad L \quad} & \textbf{LEFT_CURLY } \textit{stmt} \textbf{ RIGHT_CURLY} \\ & \xrightarrow{\quad L \quad} & \textbf{LEFT_CURLY } \textbf{RIGHT_CURLY} \end{array}$$

The application of $\textit{expr} \rightarrow \epsilon$ effectively removes the nonterminal from the derivation by replacing it with an empty string. This disappearance is important when you consider how an LL parser decides which of several productions to apply when the nonterminal being replaced has several right-hand sides. There's no problem if all of the right-hand sides start with different terminal symbols; the parser can decide on the basis of the lookahead symbol. If the current lookahead symbol matches one of the terminals, then the production whose right-hand side starts with that symbol is used. This kind of grammar—where the leftmost symbols of all the right-hand sides of any given nonterminal are different terminal symbols—is called an *S grammar*, and S grammars are among the simplest of the LL(1) grammars to recognize.

S grammars.

An ϵ production causes problems because the parser has to look beyond the current production to make its decision. The parser still decides whether or not a symbol matches an empty string by looking at the next input symbol, but it also has to look at the grammatical symbols that can follow the nonterminal that goes to ϵ . In the current example, if the next input symbol is a **NUMBER**, then the parser can apply $\textit{stmt} \rightarrow \textbf{NUMBER}$. It applies the ϵ production if the next input symbol is a **CLOSE_CURLY**. The reasoning here is that, if an empty string is matched, then the current nonterminal can effectively disappear from the derivation, and as a consequence, the next input symbol must be a symbol that can follow that nonterminal.

Since the parser can decide which production to apply in the current grammar by looking at a single lookahead character, then this is an LL(1) grammar. If a *stmt* could start with a **CLOSE_CURLY**, or if a **NUMBER** could follow a *stmt*, then the parser wouldn't be able to decide what to do, and the grammar would not be LL(1). The real situation is a little more complicated than the foregoing would indicate—I'll discuss the ins and outs of LL(1) grammars in depth in the next chapter. The current example serves to illustrate the sorts of problems that are involved, however.

3.5 The End-of-Input Marker

Another symbol that bears mentioning is the end-of-input marker, represented by \vdash . Strictly speaking, the end-of-input marker is treated as a token, and it always follows the rightmost symbol on the right-hand side of the start production—the one with the goal symbol on its left-hand side. An explicit \vdash is often omitted from a grammar, however. The \vdash is still there—it's just not shown. This omission doesn't present problems in many

The \vdash symbol.

situations. Consider the case of a Pascal program that has to end in a period. The starting production in a Pascal grammar could look like this:

$$\text{program} \rightarrow \text{definitions PERIOD}$$

and the parser would just look for the period to detect the end of the input sentence. Any input that followed the period, including the \vdash , could be ignored in this situation. The situation is complicated in C, because the start production can look like this:

$$\text{program} \rightarrow \text{definitions} \mid \epsilon$$

In this case, the parser has to look for an \vdash marker, even though there's no explicit marker in the grammar. Just mentally tack a \vdash to the end of all the right-hand sides of the start production. (Remember, ϵ is an identity element for string concatenation, so $\epsilon\vdash$ is the same thing as \vdash .)

3.6 Right-Linear Grammars

State machines and grammars.

Right-linear grammars.

Translating a right-linear grammar to a DFA.

This section moves away from a general discussion of grammars to the specifics of implementing them. Chapter Two described how state machines can be used for lexical analysis, and the notion of a regular definition, a grammatical representation of the input language, was also introduced. It turns out that all languages that can be represented as state machines can also be represented grammatically, and I'll use this similarity to demonstrate the relationships between the two systems here.

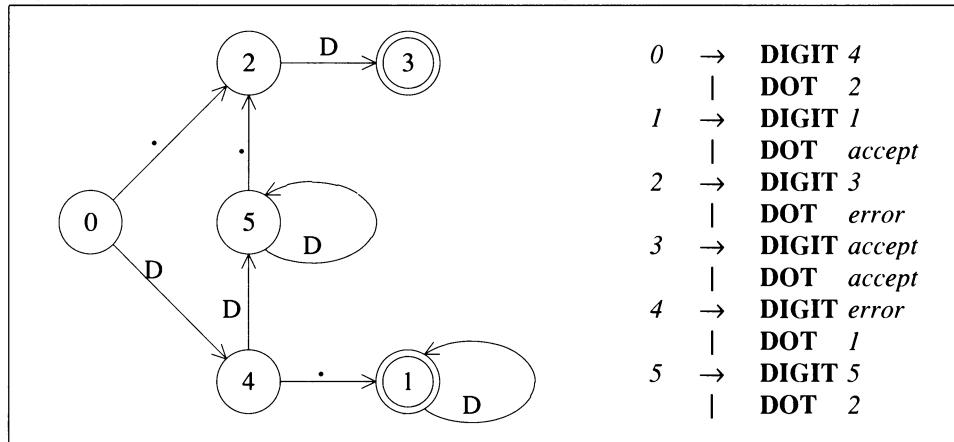
Grammars that can be translated to DFAs are called *right-linear* grammars. A grammar is right-linear if the right-hand side of each nonterminal has at most one nonterminal in it, and that nonterminal is at the far right of the right-hand side. (A *left-linear* grammar is the same, except that the nonterminal symbol must be at the far left.)

A right-linear grammar can be translated directly into a DFA if all productions in the grammar are either of the form $a \rightarrow \epsilon$ or $a \rightarrow X b$. That is, the right-hand side must either be ϵ , or it must be made up of a single terminal symbol followed by a single nonterminal symbol. To make a DFA, you must add the further restriction that, if a production has more than one right-hand side, the right-hand sides must all start with different symbols. (If they don't, you have an NFA.) The state machine is created from the grammar as follows:

- (1) The set of terminal symbols in the grammar form the DFA's input alphabet.
- (2) The set of nonterminal symbols in the grammar form the states in the DFA. The start production is the DFA start state.
- (3) If a production takes the form $a \rightarrow \epsilon$, then State a is an accepting state.
- (4) If a production takes the form $a \rightarrow X b$, then a transition is made from State a to State b on the character X .

Figure 3.2 shows both a state machine and grammar for recognizing a subset of the C floating-point constants (the regular expression $D^* \backslash . D | D \backslash . D^*$ is recognized).

A parser for this sort of grammar is, of course, trivial to implement. In fact, that's what the L^EX-generated state-machine driver is—a parser for a right-linear grammar. State transitions are made based on the current state and input character, and an action is executed when an accepting state is entered. A right-linear grammar that doesn't have the properties discussed earlier can easily be modified to have the required properties using the transformation rules discussed in Chapter Four.

Figure 3.2. State-machine and Grammar That Recognizes $D^* \cdot D \mid D \cdot D^*$ 

3.7 Lists, Recursion, and Associativity

Probably the most common grammatical structure is a list. Programs are lists of variable declarations and subroutines, subroutines are lists of statements, expressions are lists of identifiers and constants delimited by operators, and so forth. We've actually been using lists informally up until now, but it's worthwhile to look at them in depth.

3.7.1 Simple Lists

Lists are loops, and loops are implemented grammatically using recursion. You can see what's going on graphically by looking at the grammatical representations of States 1 and 5 in Figure 3.2 on page 175. The loops in these states recognize lists of digits, but you can use the same grammatical structure to recognize loops of any sort.

The recursion can be implemented in two ways. A *left-recursive* production is one where the left-hand side appears as the leftmost symbol on the right-hand side. A *right-recursive* production is the other way around, the left-hand side is duplicated at the far right. Note that a grammar can be right (or left) recursive and not be right (or left) linear because there can be more than one nonterminal on the right-hand side.

Left and right recursion.

The following grammar provides an example of a simple, left-recursive list of statements:

$$\begin{array}{lcl} \textit{stmt_list} & \rightarrow & \textit{stmt_list} \textit{ stmt} \\ & | & \textit{stmt} \end{array}$$

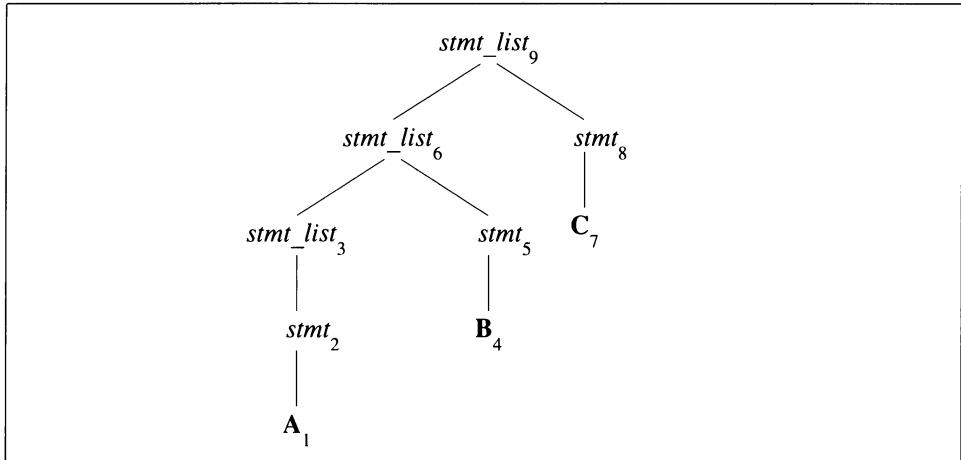
If you also supply a simple definition of *stmt* as one of three terminal symbols:

$$\textit{stmt} \rightarrow \textbf{A} \mid \textbf{B} \mid \textbf{C}$$

then the list ABC generates the parse tree in Figure 3.3.

A typical parser creates the parse tree from left to right as it reads the input, but it does a left-to-right, depth-first traversal of the parse tree as it generates code. (The tree is traversed left-to-right, but the nodes further down on the tree are visited first.) The subscripts in Figure 3.3 show the order in which nodes are visited. Assuming that the code that processes a statement is executed when the *stmt_list* nodes are traversed, the list elements are processed from left to right. That is, the elements in the list associate from left to right. This is always the case. Left-recursive productions always associate from left to right.

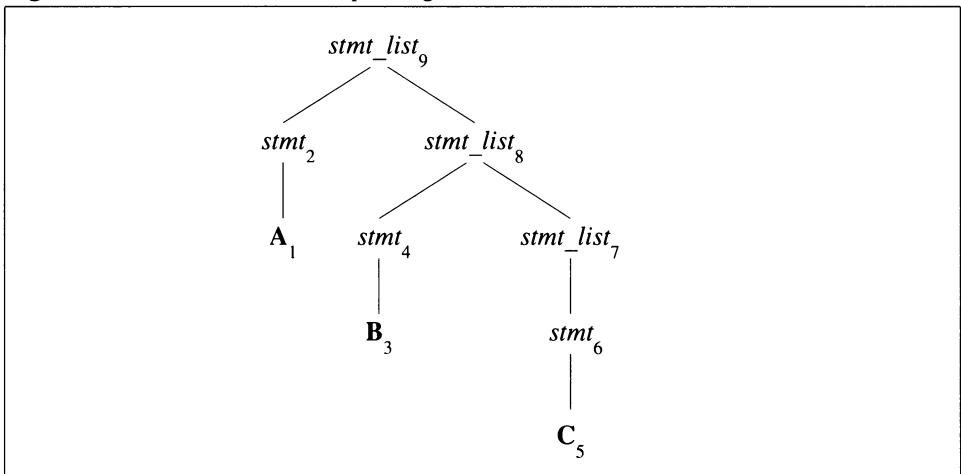
Associativity and recursion.

Figure 3.3. Parse Tree for a Simple, Left-Recursive List

Changing the grammar to a right-recursive form illustrates the change to right associativity. The new grammar looks like this:

$$\begin{aligned} \textit{stmt_list} &\rightarrow \textit{stmt stmt_list} \mid \textit{stmt} \\ \textit{stmt} &\rightarrow \textit{A} \mid \textit{B} \mid \textit{C} \end{aligned}$$

and the parse tree for ABC is shown in Figure 3.4. Again, a depth-first traversal causes the nodes to be visited in the indicated order, and the important thing is the order in which the *stmt_list* nodes are visited. The elements in this list are processed from right to left.

Figure 3.4. Parse Tree for a Simple, Right-Recursive List

Fudging left associativity with a right-recursive grammar.

Recursive-descent parsers can cheat and generate code as the tree is built from the top down. This cheating can give surprising results, because a list can be processed from left to right, even though the grammar would indicate otherwise. For example, Listing 3.1 shows the list being processed from left to right, and Listing 3.2 processes the same input from right to left. The grammar is the same, right-recursive grammar in both listings, but the position of the `process_statement()` subroutine has been changed. In Listing 3.1, `process_statement()` is called before the recursive `stmt_list()` call, so the list element is processed before the subtree is built. In the second listing, the

processing happens after the recursive call, so the subtree will have been traversed before the list element is processed.

Listing 3.1. Left Associativity with a Right-Recursive Grammar

```

1  stmt_list()
2  {
3      /* Code is generated as you create the tree, before the subtree is
4      * processed.
5      */
6
7      remember = stmt();
8
9      process_statement( remember );
10
11     if( not_at_end_of_input() )
12         stmt_list();
13 }
14
15 stmt()
16 {
17     return( read() );
18 }
```

Listing 3.2. Right Associativity with a Right-Recursive Grammar

```

1  stmt_list()
2  {
3      /* Code is generated as you create the tree, after the subtree is
4      * processed.
5      */
6
7      remember = stmt();
8
9      if( not_at_end_of_input() )
10         stmt_list();
11
12     process_stmt();
13 }
14
15 stmt()
16 {
17     return read();
18 }
```

Since you can't have left recursion in a recursive-descent parser, this technique is often useful, but it can also cause maintenance problems, because some code is generated as the tree is created, but other code must be generated as the tree is traversed. You can't apply this technique in processing expressions, for example. Consequently, it's usually best always to generate code after the call to a subroutine that visits a subtree, not before it. In Chapter Four, we'll look at a method for modifying grammars that can eliminate the left recursion but maintain the left associativity; in general, this is the best approach.

There is one maintainable way to get left associativity with a right-associative grammar, though this method can be used only occasionally. Looking again at Figure 3.4, you'll notice that the *stmt* nodes are processed from left to right, even though the

stmt_list nodes go from right to left. Consequently, if the statements can be processed in *stmt* rather than in *stmt_list*, the associativity is effectively reversed.

Strictly speaking, the grammars we've been looking at are *self* left or right recursive. *Self recursive* productions are those where the left-hand side appears at one end or the other of the right-hand side. It's also possible for a grammar to have *indirect* recursion, however. Here, one or more steps have to be added to a derivation before the recursion becomes evident. For example, the following grammar has indirect left recursion in it:

1.	$s \rightarrow a \text{ AHH}$
2.	$a \rightarrow c \text{ CHOO}$
3.	$c \rightarrow s \text{ DOO_WOP}$

because, all derivations that start with *s* ultimately end up with an *s* on the far left of the viable prefix (of the partial derivation):

$$\begin{array}{lll} s & \xrightarrow{L} & a \text{ AHH} & \text{by Production 1} \\ & \xrightarrow{L} & c \text{ CHOO AHH} & \text{by Production 2} \\ & \xrightarrow{L} & s \text{ DOO_WOP CHOO AHH} & \text{by Production 3} \end{array}$$

In general, indirect recursion is not desirable in a grammar because it masks the fact that the grammar is recursive.

3.7.2 The Number of Elements in a List

The grammar we just looked at creates a list with at least one element in it. It's possible for a list to be empty, however. The following productions recognize zero or more declarations:

$$\begin{array}{lll} \text{decl_list} & \rightarrow & \text{decl_list } \text{declaration} \\ & | & \epsilon \end{array}$$

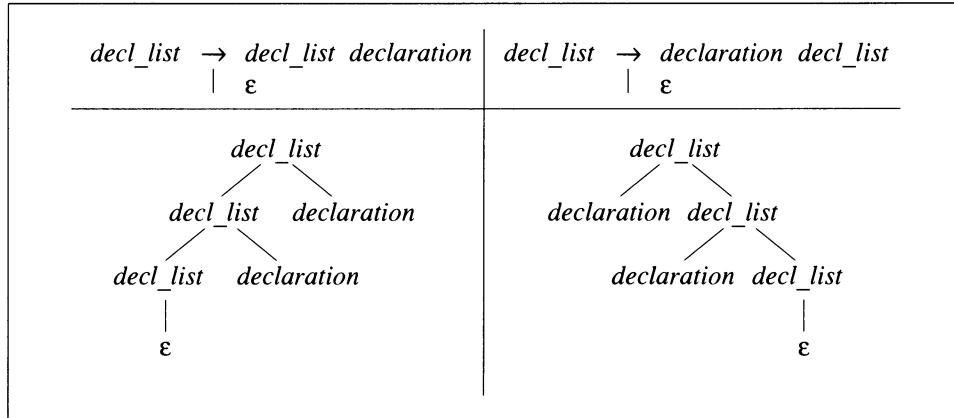
Productions executed first or last.

Since a *decl_list* can go to ϵ , the list could be empty. Parse trees for left and right recursive lists of this type are shown in Figure 3.5. Notice here, that the $\text{decl_list} \rightarrow \epsilon$ production is the first list-element that's processed in the left-recursive list, and it's the last list-element processed in the right-recursive list. The same sort of thing was also true in the nonempty lists that we looked at in the last section—the $\text{stmt_list} \rightarrow \text{stmt}$ production was processed only once in both types of lists, and it was processed first in the left-recursive list and last in the right-recursive list. This fact comes in handy when you're doing code generation, because you can use the ϵ production to trigger initializations or clean-up actions. The technique is used heavily in Chapter Six.

It's also possible to specify a limited number of repetitions, using something like the following:

$$\begin{array}{lll} \text{list} & \rightarrow & \text{element} \\ & | & \text{element element} \\ & | & \text{element element element} \end{array}$$

The method is cumbersome, but workable. Though this method is the only way to do it with a grammar, you can solve the limited-number-of-repetitions problem semantically rather than syntactically. Instead of building the counts into the grammar, you can use a more general-purpose list and have the parser keep track of how many times the list productions are applied, printing an error message if there are an incorrect number of repetitions.

Figure 3.5. Lists Terminated with ϵ Productions

3.7.3 Lists with Delimiters

Most lists must incorporate delimiters of some sort into them, and the delimiters can be handled in two ways. In C, for example, a comma separates multiple names in a declaration, as in:

```
int x, y, z;
```

Semicolons terminate elements of a list of several declarations, however. Both types of lists are illustrated in the following grammar:

<i>decl_list</i>	\rightarrow	<i>decl_list declaration</i>
		ϵ
<i>declaration</i>	\rightarrow	<i>declarator_list specifier_list SEMICOLON</i>
<i>declarator_list</i>	\rightarrow	<i>TYPE declarator</i>
		ϵ
<i>specifier_list</i>	\rightarrow	<i>specifier_list COMMA NAME</i>
		<i>NAME</i>

A *decl_list* is a list of zero or more declarations (it's left associative). The *declaration* production defines a single list element. It is made up of declarators followed by specifiers, and is terminated with a **SEMICOLON**. Since **SEMICOLON** is a terminator, it's at the far right of the production that describes the list element, rather than being in the list-production (the *decl_list*) itself. A *declarator_list* is a simple list of zero or more **TYPE** tokens. Finally, a *specifier_list* is a **COMMA**-separated list of identifiers. Here, the **COMMA**, since it is a separator, is part of the list-definition production itself.

The *specifier_list* is left recursive. A right-recursive version of the same thing could be done like this:

<i>specifier_list</i>	\rightarrow	<i>NAME COMMA specifier_list</i>
		<i>NAME</i>

The relative positions of the **COMMA** and **NAME** must also be reversed here. Common ways to do the various types of lists are all summarized in Table 3.2.

Table 3.2. List Grammars

	No Separator	
	Right associative	Left associative
At least one element	$list \rightarrow MEMBER\ list \mid MEMBER$	$list \rightarrow list\ MEMBER \mid MEMBER$
Zero elements okay	$list \rightarrow MEMBER\ list \mid \epsilon$	$list \rightarrow list\ MEMBER \mid \epsilon$
Separator Between List Elements		
	Right associative	Left associative
	$list \rightarrow MEMBER\ \text{delim}\ list \mid MEMBER$	$list \rightarrow list\ \text{delim}\ MEMBER \mid MEMBER$
Zero elements okay	$opt_list \rightarrow list \mid \epsilon$ $list \rightarrow MEMBER\ \text{delim}\ list \mid MEMBER$	$opt_list \rightarrow list \mid \epsilon$ $list \rightarrow list\ \text{delim}\ MEMBER \mid MEMBER$

A MEMBER is a list element; it can be a terminal, a nonterminal, or a collection of terminals and nonterminals. If you want the list to be a list of terminated objects such as semicolon-terminated declarations, MEMBER should take the form: MEMBER $\rightarrow \alpha\ \text{TERMINATOR}$, where α is a collection of one or more terminal or nonterminal symbols.

3.8 Expressions

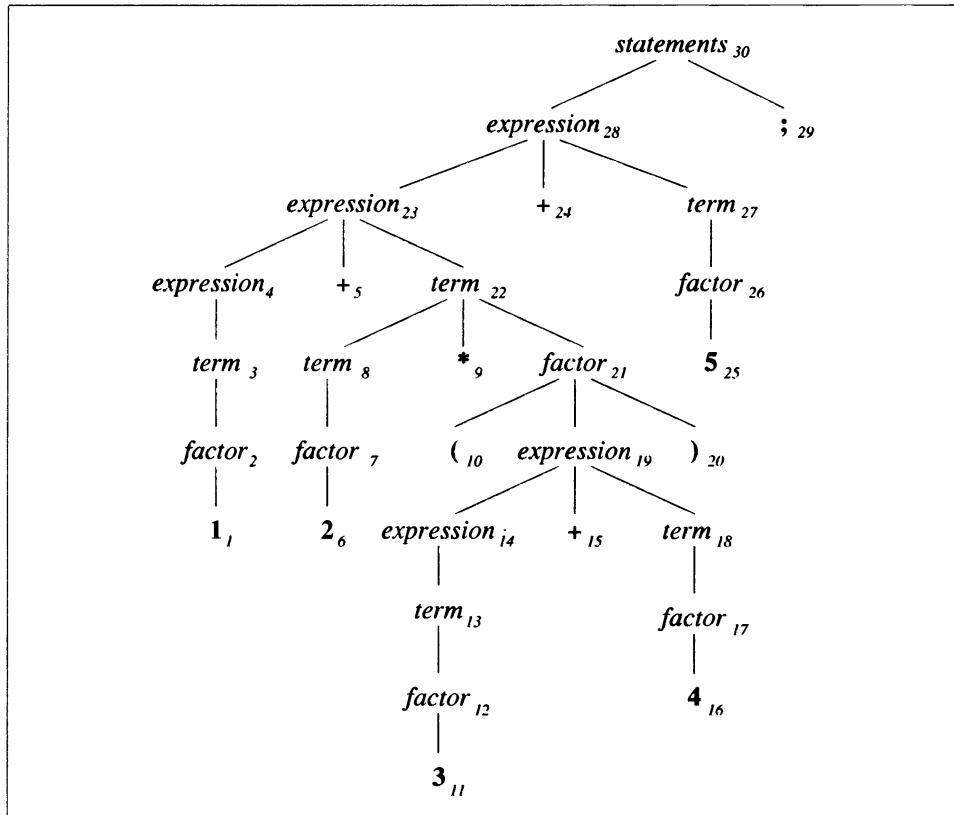
One special form of a list is an expression. For example $1+2+3$ is a plus-sign-delimited list of numbers. Table 3.3 shows a grammar that recognizes a list of one or more statements, each of which is an arithmetic expression followed by a semicolon. A *stmt* is made up of a series of semicolon-delimited expressions (*exprs*), each a series of numbers separated either by asterisks (for multiplication) or plus signs (for addition). This grammar has several properties that are worth discussing in depth.

Table 3.3. A Simple Expression Grammar

1.	<i>stmt</i>	\rightarrow	<i>expr</i> ;
2.			<i>expr</i> ; <i>stmt</i>
3.	<i>expr</i>	\rightarrow	<i>expr</i> + <i>term</i>
4.			<i>term</i>
5.	<i>term</i>	\rightarrow	<i>term</i> * <i>factor</i>
6.			<i>factor</i>
7.	<i>factor</i>	\rightarrow	number
8.			(<i>expr</i>)

Depth-first traversal.

Figure 3.6 shows the parse tree generated by the foregoing grammar when the input statement $1+2*(3+4)+5;$ is processed. The code-generation pass does a *depth-first* traversal of the tree—the deeper nodes are always processed first. The subscripts in the figure show the order in which the nodes are visited. Assuming that code is generated only after an entire right-hand side has been processed, the depth-first traversal forces the $3+4$ to be done first (because it's deepest in the tree and the *expr* associated with the subexpression $3+4$ is processed before any of the other *exprs*); then, moving up the tree, the result of the previous addition is multiplied by the 2; then 1 is added to the accumulated subexpression; and, finally, the 5 is added into the total.

Figure 3.6. A Parse Tree for $1+2*(3+4)+5;$ 

Note that the order of evaluation is what you would expect—assuming that multiplication is higher precedence than addition, that addition associates from left to right, and that parenthesized expressions are evaluated first. Since the parse tree in Figure 3.7 is the only possible one, given the grammar used, this means that associativity and precedence are actually side effects of the grammatical structure. In particular, subtrees that are lower down on the parse tree are always processed first, and positioning a production further down in the grammar guarantees a lower relative position in the parse tree. Higher-precedence operators should then be placed lower down in the grammar, as is the case in the grammar used here. (The only way to get to a *term* is through an *expr*, so multiplication is always further down the tree than addition.) Parentheses, since they are at the bottom of the grammar, have higher precedence than any operator. The associativity is still determined by left or right recursion, as is the case in a simple list.

Unary operators tend to be lower in the grammar because they are usually higher precedence than the binary operators. You can add unary minus to the previous grammar by modifying the factor rules as follows:

Order of evaluation as controlled by grammar.

Unary operators.

1.	<i>stmt</i>	\rightarrow	<i>expr</i> ;
2.			<i>expr</i> ; <i>stmt</i>
3.	<i>expr</i>	\rightarrow	<i>expr</i> + <i>term</i>
4.			<i>term</i>
5.	<i>term</i>	\rightarrow	<i>term</i> * <i>unop</i>
6.			<i>unop</i>
7.	<i>unop</i>	\rightarrow	- <i>factor</i>
8.			<i>factor</i>
9.	<i>factor</i>	\rightarrow	number
10.			(<i>expr</i>)

The placement of the new productions insure that unary minus can be applied to parenthesized subexpressions as well as single operands.

Unary operators can sometimes be treated as lists. For example, the C pointer-dereference operators (*), can pile up to the left of the operand, like this: ***p. You can introduce a star into the previous grammar by adding the following list production:

$$\text{factor} \rightarrow * \text{ factor}$$

Note that the right recursion correctly forces right-to-left associativity for this operator.

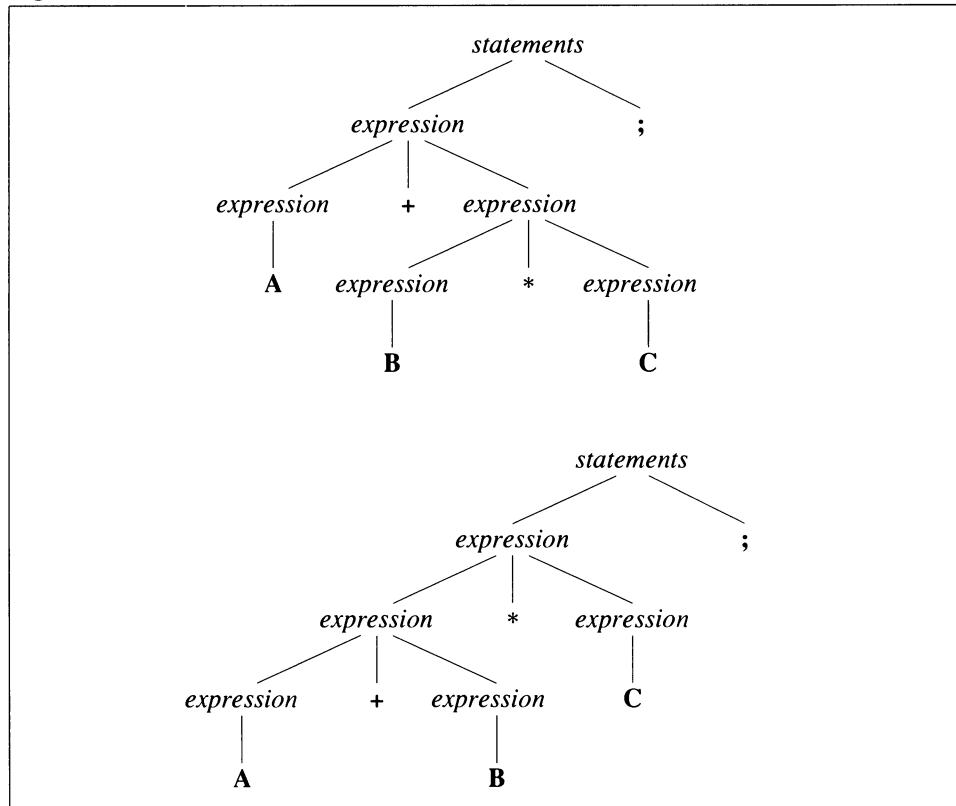
3.9 Ambiguous Grammars

The expression grammar just discussed is a *unambiguous* grammar because only one possible parse tree can be created from any given input stream. The same parse tree is generated, regardless of the derivation used. It's possible to write an *ambiguous* grammar, however. For example, expressions could be represented as follows:

$$\begin{array}{lcl} \text{statement} & \rightarrow & \text{expr} ; \\ \text{expr} & \rightarrow & \text{expr} + \text{expr} \\ & | & \text{expr} * \text{expr} \\ & | & (\text{expr}) \\ & | & \text{number} \end{array}$$

A grammar is ambiguous when the same nonterminal appears twice on a right-hand side (because the order in which the nonterminals are evaluated is dependent on the derivation). Two of the possible parse trees that can be generated for A+B*C are shown in Figure 3.7.

Ambiguous grammars are generally to be avoided exactly because they're ambiguous. That is, because there are two possible parse trees, the expression can be evaluated in two different ways, and there's no way to predict which one of these ways will be used from the grammar itself. Precedence and associativity can't be controlled by the grammar alone. Ambiguous grammars do tend to be smaller, however, and easier to read. So, parser-generation programs like **yacc** and **occs** generally provide mechanisms for using them. These programs can force specific associativity or precedence by controlling the parse in predetermined ways when an ambiguous production is encountered. The technique is discussed in depth in Chapter Five.

Figure 3.7. Parse Trees Derived from an Ambiguous Grammar

3.10 Syntax-Directed Translation

I mentioned earlier that if the syntactic scope of a language was sufficiently limited, you could derive semantic information from a formal grammar—from the position of particular symbols in the input sentence, and by extension in the grammar itself. This is how compilers, in fact, generate code—by executing code-generation actions at times that are controlled by the positions of various strings in the input sentence. We did this on an ad-hoc basis in Chapter One, first by building a parser for a grammar, and then by adding code-generation actions to the parser. The process is discussed in greater depth in this section.

3.10.1 Augmented Grammars

In an augmented grammar, code-generation actions are placed in the grammar itself, and the relative position of the action in the grammar determines when it is executed. For example, a production like this:

$$\text{expr}' \rightarrow + \text{term} \{ \text{op} ('+') ; \} \text{expr}'$$

could be coded in a recursive-descent parser as follows:

Augmentations.

```

expr_prime()
{
    if( match( PLUS ) )
    {
        term();
        op('+');
        expr_prime();
    }
}

```

Generating code to evaluate expressions.

I'll demonstrate how an augmented grammar works by generating code to evaluate expressions. You need two subroutines for this purpose, both of which generate code. In addition, between the two of them, they manage a local stack of temporary-variable names. An expression of the form $1+2+3$ generates the following code:

```

t0 = 1;
t1 = 2;
t0 += t1;
t1 = 3;
t0 += t1;

```

Anonymous temporaries.

The variables $t0$ and $t1$ are known as *anonymous temporaries* (or just plain *temporaries*) because the compiler makes up the names, not the programmer. The actual code is generated with the following sequence of actions (the numbers identify individual actions, several of which are repeated):

- (1) Read the 1.
- (2) Get a temporary-variable name ($t0$), and output $t0 = 1$.
- (3) Push the temporary-variable name onto a local stack.
- (4) Read the "+" and remember it.
- (1) Read the 2.
- (2) Get a temporary-variable name ($t1$), and output $t1 = 2$.
- (3) Push the temporary-variable name onto a local stack.
- (5) Pop two temporary-variable names off the stack and generate code to add them together: $t0 += t1$ [you're adding because you remembered it in (4)].
- (6) Push the name of the temporary that holds the result of the previous operation and discard the second one.
- (4) Read the "+" and remember it.
- (1) Read the 3.
- (2) Get a temporary-variable name ($t1$), and output $t1 = 3$.
- (3) Push the temporary name onto a local stack.
- (5) Pop two temporary-variable names off the stack and generate code to add them together: $t0 += t1$.
- (6) Push the name of the temporary that holds the result of the previous operation and discard the second one.

When you're done, the name of the anonymous temporary that holds the evaluated expression is at the top of the stack. Many of the previous actions are identical—the time that the action is executed is the important factor here.

The subroutines that do the work are in Listing 3.3. Temporary-variable names are allocated and freed with the `newname()` and `freename()` calls on lines nine and ten using the method described in Chapter One. You allocate a name by popping it off a stack (declared on line six), and free the name by pushing it back. I've removed the error-detection code to simplify things.

The rest of Listing 3.3 shows the code-generation actions. The `create_tmp()` subroutine (on line 14) does two things, it generates the code that copies the current lexeme

Listing 3.3. *expr.y*—Action Code for the Augmented Grammar

```

1 #include <tools/stack.h>                                /* Described in Appendix A. */
2
3 stack_dcl(Temporaries, char*, 128);    /* Temporaries: stack of 128 char ptrs. */
4 /*-----*/
5
6 char *Names[] = { "t0", "t1", "t2", "t3", "t4", "t5", "t6", "t7" };
7 char **Namep = Names;
8
9 char *newname()           { return( *Namep++ ); }
10 void freename(s) char *s; { *--Namep = s; }
11
12 /*-----*/
13
14 create_tmp( str )      /* Create a temporary holding str and push its name. */
15 char *str;
16 {
17     char *name = getname();
18     yy_code("%s = %s;\n", name, str );
19     push( Temporaries, name );
20 }
21
22 op( what )
23 int what;
24 {
25     char *left, *right;
26
27     right = pop( Temporaries );
28     left = tos( Temporaries );
29
30     yy_code("%s %c= %s;\n", left, what, right );
31     freename(right);
32 }
```

into an anonymous temporary, and it pushes the name of that temporary onto the Temporaries stack. The `op()` subroutine on line 22 generates code to do the actual operation. It takes two temporary-variable names from the Temporaries and generates the code necessary to perform the required operation. Note that only one of the temporaries is actually popped. The other remains on the stack because it holds the result of the generated operation. The popped name is recycled for future use on line 31.

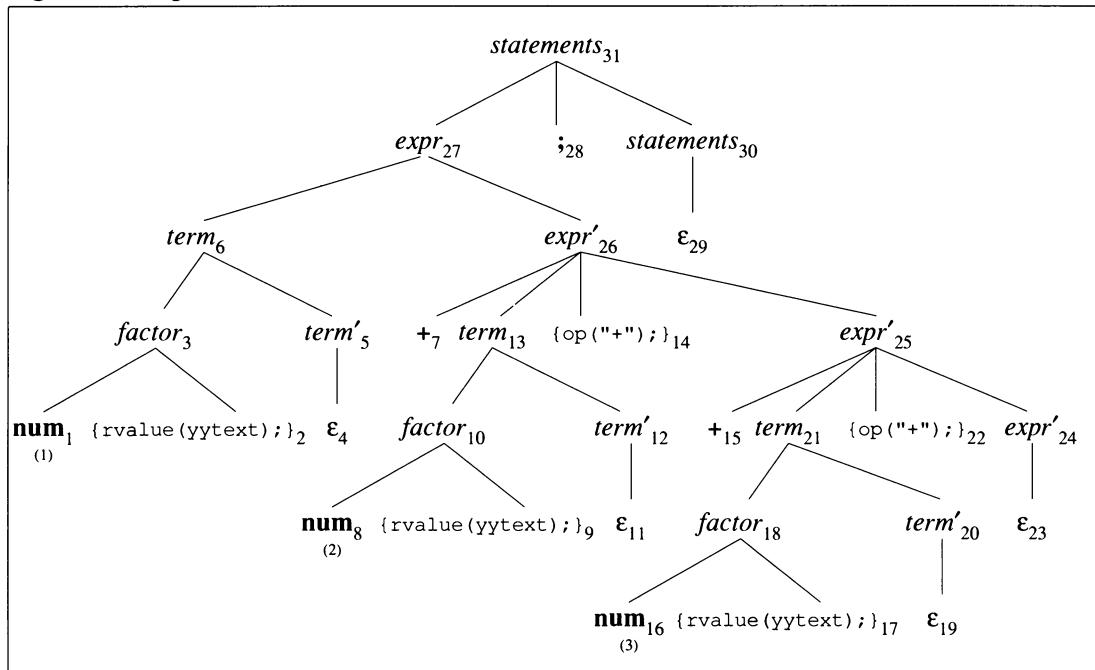
Now, armed with the proper tools, you can design the parser. The augmented grammar in Table 3.4 performs the actions described earlier in the correct order. The “remember it” action is implicit in the grammar itself (because, to get to the `op()` call in Production 4, for example, you must have passed a previous `+`). The order of evaluation is controlled, as usual, by the structure of the parse tree and the fact that you’re doing a bottom-up traversal. Action symbols in the grammar are treated like other terminal symbols—they’re just put onto the parse tree in the appropriate places and executed when that node is traversed. A parse tree for $1+2*3$ is shown in Figure 3.8. The subscripts indicate the order in which nodes are visited.

Designing the parser.

An exercise is in order at this juncture. Take a pencil and piece of paper and draw both the parse tree and the contents of the Temporaries stack (which holds the temporary-variable names) as the expression $1+2*3$ is parsed. Note how the Temporaries stack keeps track of the name of the variable that holds the 1 until after the higher-precedence multiplication operator is processed. The 2 is stored in the same way until the parenthesized subexpression can be handled. Also, note how the name at top of

Table 3.4. An Augmented Expression Grammar

1.	<i>stmt</i>	$\rightarrow \epsilon$
2.		$ \quad expr ; stmt$
3.	<i>expr</i>	$\rightarrow term \ expr'$
4.	<i>expr'</i>	$\rightarrow + \ term \ \{ op ('+') ; \} \ expr'$
5.		$ \quad \epsilon$
6.	<i>term</i>	$\rightarrow factor \ term'$
7.	<i>term'</i>	$\rightarrow * \ factor \ \{ op ('*') ; \} \ term'$
8.		$ \quad \epsilon$
9.	<i>factor</i>	$\rightarrow \text{number_or_id} \ \{ \text{create_tmp}(yytext) ; \}$
10.		$ \quad (\ expr)$

Figure 3.8. Augmented Parse Tree for $1+2+3$ 

stack always holds the value of the most recently processed subexpression.

3.10.2 Attributed Grammars

Attributes.

So far, the concept of an *attribute* has been used in a limited, informal way. An attribute is a piece of information that is associated with a grammatical symbol. Tokens all have at least one attribute—the associated lexeme. They can have other attributes as well. For example a **NUMBER** token could have an integer attribute that was the number represented by the lexeme. A **NAME** token could have an attribute that was a pointer to a symbol-table entry for that name. Attributes of this type are usually represented with subscripts. For example, the two attributes attached to a **NUMBER** could be represented as follows:

NUMBER_(“156”, 156)

The first attribute is the lexeme, the second is the numeric value of the lexeme.

Nonterminal symbols can have attributes as well. Take, for example, a recursive-descent parser. You can look at the subroutines that implement nonterminals as symbols that represent the nonterminal, not the other way around. Looked at in this way, a subroutine's argument or return value represents a quantum of information that is attached to the grammatical symbol—to the nonterminal that the subroutine is implementing, so the subroutine arguments and return values can be viewed as the attributes of the associated nonterminal symbol.

Attributes for nonterminals.

There are two types of attributes that are attached to nonterminals, and these two types correspond to the two methods used to pass information around in the two parsers presented at the end of Chapter One. An *inherited* attribute is passed down the parse tree, from a parent to a child, so in a recursive-descent compiler, inherited attributes are subroutine arguments. A *synthesized* attribute goes in the other direction, up the tree from a child to a parent. In a recursive-descent compiler, synthesized attributes are return values.¹ The best way to keep track of which term is which is to think of the actual English usage of the two words. Children “inherit” something from their parents, not the other way around. So, an “inherited” attribute is one that’s passed from the parent to the child in the parse tree. The code-generation in the previous section was done without the benefit of attributes. That is, all information was passed between code-generation actions using global variables rather than attaching the information to the grammatical elements themselves. Drawing an analogy to recursive-descent parsers, it’s as if you didn’t use subroutine arguments or return values anywhere in the parser.

Inherited and synthesized attributes.

The attribute mechanism can be extended to grammars in general and is very useful in designing a compiler. The attributes can help you specify code-generation actions in the grammar in greater detail than would be possible with an augmented grammar alone. A grammar to which attributes have been added in this way is called an *attributed grammar*. In a typical compiler design, you put attributes into the grammar in a manner similar to the augmentations used in the previous sections. That is, the first step in the design process is the creation of an augmented and *attributed* grammar. It’s easier to understand the concepts, however, if you start with working code, so I’ll demonstrate the process by going backwards from a working parser to an attributed, augmented grammar. The naive recursive-descent parser in Listing 3.4 implements the following grammar:

Attributed grammar.

1.	<i>stmt</i>	\rightarrow	ϵ
2.			<i>expr ; stmt</i>
3.	<i>expr</i>	\rightarrow	<i>term expr'</i>
4.	<i>expr'</i>	\rightarrow	$+ \text{term } \textit{expr}'$
5.			ϵ
6.	<i>term</i>	\rightarrow	<i>factor term'</i>
7.	<i>term'</i>	\rightarrow	$* \text{factor } \textit{term}'$
8.			ϵ
9.	<i>factor</i>	\rightarrow	number
10.			(<i>expr</i>)

This parser is essentially the naive parser developed in Chapter One, where you’ll find the `match()` and `advance()` subroutines that comprise the lexical analyzer, `newname()` and `freename()` were discussed earlier. The parser uses inherited

1. In ALGOL-like languages such as Pascal, an inherited attribute is a subroutine argument that’s passed by value; a synthesized attribute is either passed by reference or is a function’s return value.

attributes, which are passed down the parse tree from parent to child as subroutine arguments.

Listing 3.4. *naive.c*—Code Generation with Synthesized Attributes

```

1 void factor ( char *t );                                /* Prototypes to avoid */
2 void term   ( char *t );                                /* forward references */
3 void expr   ( char *t );                                /* in the code. */
4
5 stmt()          /*  stmt -> expr SEMI  |  expr SEMI stmt    */
6 {
7     char *t;
8
9     while( !match(EOI) )
10    {
11        expr( t = newname() );
12        freename( t );
13
14        if( match( SEMI ) )
15            advance();
16    }
17 }
18 -----
19 void expr( t )                                     /* expr  -> term expr' */
20 char *t;
21 {
22     term( t );
23     expr_prime( t );
24 }
25 -----
26 expr_prime( t )                                    /* expr' -> PLUS term expr' | epsilon */
27 char *t;
28 {
29     char *t2;
30
31     if( match( PLUS ) )
32     {
33         advance();
34
35         term( t2 = newname() );
36
37         printf(" %s += %s\n", t, t2 );
38         freename( t2 );
39
40         expr_prime( t );
41     }
42 }
43 -----
44 void term( t )                                     /* term  -> factor term' */
45 char *t;
46 {
47     factor   ( t );
48     term_prime ( t );
49 }
50 -----
51 term_prime( t )                                    /* term' -> TIMES factor term' | epsilon */
52 char *t;
53 {
54     char *t2 ;
55 }
```



Listing 3.4. continued...

```

56     if( match( TIMES ) )
57     {
58         advance();
59
60         factor( t2 = newname() );
61
62         printf("    %s *= %s\n", t, t2 );
63         freename( t2 );
64
65         term_prime( t );
66     }
67 }
68 -----*/
69 void factor( t )           /* factor -> NUMBER_OR_ID | LP expr RP */
70 char *t;
71 {
72     if( match(NUMBER_OR_ID) )
73     {
74         printf("    %s = %0.*s\n", t, yyleng, yytext );
75         advance();
76     }
77     else if( match(LP) )
78     {
79         advance();
80
81         expr( t );
82
83         if( match(RP) )
84             advance();
85     }
86 }
```

Now, consider the flow of attributes (the temporary-variable names) through the parse tree as the expression $1+2*3$ is processed. The parser generates the following output:

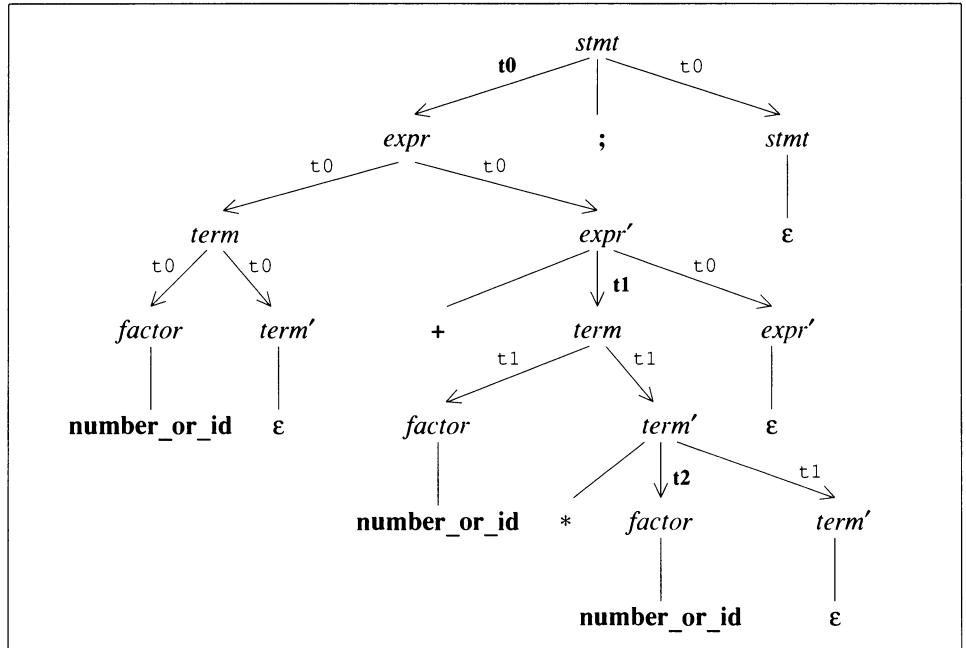
```
t0 = 1
t1 = 2
t2 = 3
t1 *= t2
t0 += t1
```

and the associated parse tree is shown in Figure 3.9. The flow of attributes is shown in this figure by labeling edges in the graph with the attribute's value. That is, the fact that *stmt()* passes *t0* to *expr()* is indicated by labeling the edge between the *stmt* and *expr* nodes with a *t0*. The attribute is in boldface when the calling function creates that attribute with a *newname()* call; otherwise, the attribute came into a function as an argument and was passed to a child function. You should trace through the first few subroutine calls in the parse to see what's happening. Remember, each child node in the tree represents a subroutine call made by the parent, and the attributes are subroutine arguments, so they flow down the tree from the parent to the child.

You can represent the flow of attributes by borrowing the notation used for arguments in a programming language. Each attribute has a name just like subroutine arguments have names. All attributes that are passed into a production are listed next to the name on the left-hand side of that production, like a formal argument list in a subroutine declaration. Those attributes can then be referenced by name by the symbols on the

Borrow attribute notation
from subroutine-calling
conventions.

Figure 3.9. The Flow of Inherited Attributes Through a Parse Tree for $1+2*3$



right-hand side of the production, as if they were used in the subroutine. For example, say that an attribute, t , represents a temporary variable name, and it is attached to an $expr$; it is represented like this in the subroutine representing the $expr$:

```
expr( t )
char *t;
{
    ...
}
```

and like this in the grammar:

$$expr_{(t)} \rightarrow \dots$$

A t attribute would also be attached to every $expr$ on a right-hand side, as if it were the argument to a recursive subroutine call:

$$stmt \rightarrow expr_{(t)} ; stmt$$

If the same attribute appears on both the left- and right-hand sides, then that attribute comes into the subroutine as an argument, and is passed down to a child subroutine, in turn. An attribute can also be used internally in an action, as if it were a variable.

Attributes can also be created within an action in a manner analogous to local variables, and that new attribute can be, in turn, passed to another grammatical symbol (to a subroutine) as an attribute (as an argument). Note that a leftmost derivation requires the flow of attributes to be from left-to-right across the production, just like the flow of control in a recursive-descent parser that implements the grammar. So, in an LL parser (which uses a leftmost derivation), inherited attributes can be passed from a left-hand side to a symbol on the right-hand side, and an attribute can be passed from a code-generation action to a symbol to its right. Synthesized attributes are used differently, and are discussed in depth in Chapter Five along with the bottom-up parsers that use this type of attribute.

An attributed version of our expression grammar is shown in Table 3.5. The attribute names are taken from the previous recursive-descent compiler. An attribute attached to a left-hand side is treated like a formal argument in a subroutine declaration—it's the attribute that is passed to the current production from its parent in the tree. An attribute on the right-hand side is treated like an argument in a subroutine call; it is passed down the tree to the indicated production.

Table 3.5. An Attributed Grammar

1.	<i>stmt</i>	\rightarrow	ϵ
2.			<i>expr</i> _(t) ; <i>stmt</i>
3.	<i>expr</i> _(t)	\rightarrow	<i>term</i> _(t) <i>expr'</i> _(t)
4.	<i>expr'</i> _(t)	\rightarrow	+ <i>term</i> _(t2) <i>expr'</i> _(t)
5.			ϵ
6.	<i>term</i> _(t)	\rightarrow	<i>factor</i> <i>term'</i>
7.	<i>term'</i> _(t)	\rightarrow	* <i>factor</i> _(t2) <i>term'</i> _(t)
8.			ϵ
9.	<i>factor</i> _(t)	\rightarrow	number_or_id
10.			(<i>expr</i> _(t))

The attributes by themselves are not much use. You need to augment the grammar to show how the attributes are manipulated. Remember that the flow of control in a top-down parser is from left to right across the grammar, so an action can affect only those symbols to its right and can use only those symbols that are initialized to its left (or which appear as an attribute of the left-hand side). An augmented, attributed grammar is shown in Table 3.6. The code-generation actions are taken from the previous recursive-descent parser.

Augmented, attributed grammar.

Table 3.6. An Augmented, Attributed Grammar

1.	<i>stmt</i>	\rightarrow	ϵ
2.			{t=newname();} <i>expr</i> _(t) ; <i>stmt</i>
3.	<i>expr</i> _(t)	\rightarrow	<i>term</i> _(t) <i>expr'</i> _(t)
4.	<i>expr'</i> _(t)	\rightarrow	+ {t2=newname();} <i>term</i> _(t2) {printf("%s+=%s\n", t, t2); freename(t2);} <i>expr'</i> _(t)
5.			ϵ
6.	<i>term</i> _(t)	\rightarrow	<i>factor</i> <i>term'</i>
7.	<i>term'</i> _(t)	\rightarrow	* {t2=newname();} <i>factor</i> _(t2) {printf("%s+=%s\n", t, t2); freename(t2);} <i>term'</i> _(t)
8.			ϵ
9.	<i>factor</i> _(t)	\rightarrow	number_or_id { printf("%s=%0.*s\n", t, yylen, yytext); }
10.			(<i>expr</i> _(t))

As I said at the beginning of this section, I've demonstrated the process in a topsyturvy fashion. Generally, you would create the augmented, attributed grammar first, and then use that grammar as a specification when coding. The compiler-generation tools developed in the next two chapters take augmented, attributed grammars as their input, and output parsers for these grammars, thereby eliminating the necessity of coding a parser at all. You still have to write the grammar, however, and figure out how to generate code by adding actions to it in appropriate places.

One final note: It's possible for a single grammar to use both inherited and synthesized attributes, but it's not a good idea. In a recursive-descent parser—the only kind

of parser in which combined attribute types are really practical—maintenance becomes a problem because it's more difficult to follow the flow of data as the parser executes. Other types of parsers, such as the one-pass, table driven ones discussed in the next two chapters, can handle only one type of attribute.

3.11 Representing Generic Grammars

Generic symbols.

As we discuss grammars in the remainder of this book, it will be handy, occasionally, to discuss generic productions. In these discussions, when I don't say otherwise the term *symbol*, without the word “terminal” or “nonterminal,” is used for when something can be either a terminal or nonterminal. As usual, italics are used for nonterminals and bold-face is used for terminals, but I'll use upper-case letters that look like this:

$$\mathcal{A} \mathcal{B} \mathcal{C} \mathcal{D} \mathcal{E} \mathcal{F} \mathcal{G} \mathcal{H} \mathcal{I} \mathcal{J} \mathcal{K} \mathcal{L} \mathcal{M} \mathcal{N} \mathcal{O} \mathcal{P} \mathcal{Q} \mathcal{R} \mathcal{S} \mathcal{T} \mathcal{U} \mathcal{V} \mathcal{W} \mathcal{X} \mathcal{Y} \mathcal{Z}$$

to represent an unspecified symbol. A production like this:

$$s \rightarrow \mathcal{A} \ B$$

has two symbols on its right-hand side, and these can be either terminals or nonterminals. The following production has three symbols on its right-hand side: the first is a nonterminal, the second is a terminal, and the third can be either a terminal or nonterminal:

$$s \rightarrow \alpha \ T \ \mathcal{A}$$

Greek letters represent sequence of symbols.

To confuse matters further, I'll use Greek letters to represent a sequence of symbols. For example, a generic production can be written like this:

$$s \rightarrow \alpha$$

s is a nonterminal because it's in italics, and α represents an arbitrary sequence of terminals and nonterminals. Unless I say otherwise, a Greek letter represents zero or more symbols, so something like this:

$$s \rightarrow \alpha A \beta$$

represents all productions that have s on their left-hand side and an A somewhere on their right-hand side. The A can be preceded by zero or more symbols (terminals or nonterminals), and it can be followed by zero or more symbols. A sequence of symbols represented by a Greek letter can go to ϵ . That is, it can be replaced in the equation by an empty string, effectively disappearing from the production. For example, the production

$$s \rightarrow A \alpha B$$

Multiple right-hand sides,
 α_n .

has to start with an A and end with a B , but it can have any number of terminals and nonterminals (including zero) in between.

If a generic production has several right-hand sides, I'll represent these with subscripts. For example, in the following generic grammar, the nonterminal s has n different right-hand sides, one of which could be ϵ :

$$\begin{array}{rcl} s & \rightarrow & \alpha_1 \\ & | & \alpha_2 \\ & \dots & \\ & | & \alpha_n \end{array}$$

Multiple tokens, T_n .

By the same token (so to speak), the following represents n different terminal symbols: $T_1 \ T_2 \dots T_n$.

All the foregoing may seem abstruse, but it will sink in once you've used it a few times. I'll call out exactly what's going on if the generic production gets too complicated.

3.12 Exercises

- 3.1. Show a step-by-step leftmost derivation of the expression

$1 + 2 * ((3+4) + 5) + 6;$

using the following grammar:

1.	<i>statements</i>	\rightarrow	ϵ
2.		$ $	<i>expr ; statements</i>
3.	<i>expr</i>	\rightarrow	<i>term expr'</i>
4.	<i>expr'</i>	\rightarrow	$+ \text{ term } \text{expr}'$
5.		$ $	ϵ
6.	<i>term</i>	\rightarrow	<i>factor term'</i>
7.	<i>term'</i>	\rightarrow	$* \text{ factor } \text{term}'$
8.		$ $	ϵ
9.	<i>factor</i>	\rightarrow	number
10.		$ $	(expr)

- 3.2. Build a recursive-descent parser that implements the augmented grammar in Table 3.6 on page 191.

- 3.3. The following regular expression recognizes a series of binary digits, with the pattern **000** used to terminate the series:

$(0|1)^*000$

Translate that expression into a right-linear grammar.

- 3.4. Assuming that each character in the ASCII character set forms a token, create a grammar that recognizes all C integer and character constants. All of the following should be recognized by your grammar:

`0abcd 0123 45 'a' '\t' '\x0a' '\123'`

- 3.5. In LISP, a simple expression is formed as follows: $(+ a b)$ represents $a+b$, $(+ a (* b c))$ represents $a+b*c$, and so forth. Write a grammar that recognizes all such LISP expressions, with an arbitrary level of parentheses nesting. The $+$, $-$, $*$, and $/$ operators must be supported, $+$ and $-$ should be lower precedence than $*$, and $/$.

- 3.6. Write an expression grammar that supports variable names, an assignment operator, a plus operator and array references. Array references are done in a FORTRAN-like fashion: a three-dimensional array element can be accessed using the notation $a[x, y, z]$. x is the minor axis, so the list of array indexes should be processed from right to left. Similarly an expression like $a = b[2] = c[1,2]$ should be recognized, and the assignment operator should associate right to left. Addition associates left to right, however. The order of precedence should be brackets (highest precedence), then addition, then assignment.

- 3.7. Write a grammar that recognizes Pascal subroutine declarations. Draw the parse tree that results from a top-down parse of the following declaration:

```
function subr( var arg1: string; x: integer; s: string):
    integer;
```

- 3.8. Write a grammar for 8086 assembly language (or whatever assembly language that you know).
- 3.9. Write a grammar that describes the input language used by your calculator. Every key on the calculator represents a token in the language.
- 3.10. Listing 3.5 shows a specification for a DFA transition matrix. The state statements define the rows of the matrix; each **goto** statement controls the contents of a column in that row. The accept statements say that the current state is an accepting state, and the code that follows is to be executed when the accepting action occurs. The outer braces are part of the accept statement itself, but the included code must also be able to have braces in it. Unspecified transitions are filled with the number specified in the **error** statement.

Listing 3.5. A State Machine Specification

```

1  rows 4;          /* Number of rows      in the table.           */
2  columns 128;     /* Number of columns in the table.          */
3  error -1;        /* Use this value for all unspecified table elements. */
4
5  state 0
6  {
7      goto 0 on 'a';
8      goto 1 on 'b';
9      goto 2 on 'c';
10     goto 2 on 100;    /* = 0x64 = 'd'. */
11     loop on 0x65;    /* 0x65 = 'e', same as "goto <current state>." */
12     accept           /* Define an accepting action for this state. */
13     {
14         execute_this_code();
15     }
16 }
17
18 state 1 { goto 2 on b; }           /* Not accepting. */
19 state 2 { goto 3 on c; }           /* Not accepting. */
20 state 3 { accept{ hi_there("BOO!"); } }
```

- (a) Write a grammar for the language just described.
- (b) Augment the grammar with sufficient code-generation actions and attributes that a parser could be coded that reads in a state-machine specification, and which outputs C source code, which, when compiled, implements those tables. Two array declarations and a subroutine should be output. The first array should be the transition matrix itself, the second array should be indexed by state number and should evaluate to 1 if that state is an accepting state. The subroutine should contain a **switch** statement which, when passed the state number of an accepting state, executes the associated code.
- (c) Implement the grammar developed in (b).
- 3.11. Modify your solution to the foregoing exercise so that you can give symbolic names to a state and then use those symbolic names in **goto** statements. Names are declared implicitly by using them instead of a number in a state statement. The compiler must assign state numbers in this case, and you should make sure that a state number supplied by the compiler doesn't conflict with one supplied in a state statement. Forward references should be permitted.

Top-Down Parsing

This chapter develops the idea of recursive-descent parsing by discussing ways to do the same thing in a more maintainable, table-driven fashion. The chapter includes a description of top-down parsing techniques and an in-depth discussion of LL grammars, including techniques for modifying grammars to be LL(1). A *yacc*-like utility called **LLama**—which translates an augmented, attributed LL(1) grammar into a parser—is constructed. If you intend to read the implementation parts of this chapter, you should read Appendix E, which contains a user’s manual for **LLama** and **occs**, before continuing.

One thing I’m not doing here is presenting an extended example of how to use **LLama**. I’ve left this out because **LLama** is, in itself, not nearly as useful a tool as is **occs**. The main reason this chapter is here, in fact, is to present several concepts and procedures that are used later to construct **occs** (but in a somewhat simpler context than **occs** itself). That is, **LLama**, is just a step on the way to **occs**. I do present a short example of using **LLama** in that I use **LLama** to rewrite its own parser. Also, a significant part of the code used by **LLama** is also used by **occs**, and the common code is presented only in the current chapter.

4.1 Push-Down Automata*

We saw in Chapter Three that certain right-linear grammars can be converted directly to DFAs. If you can represent a programming language in this way, you can implement a parser for it as a state machine with a code-generation action associated with each state. Consider the following simple grammar that recognizes expressions of the form **number + number**.

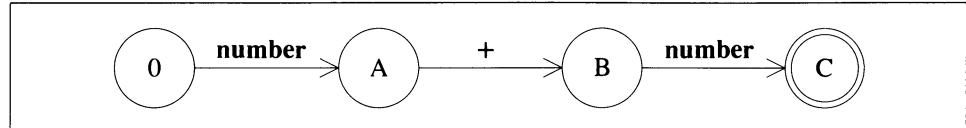
Parsing with a state machine.

* As in Chapter Two, asterisks mark sections containing theoretical material.

$$\begin{aligned}
 E &\rightarrow \text{number } A \\
 A &\rightarrow + B \\
 B &\rightarrow \text{number } C \\
 C &\rightarrow \epsilon
 \end{aligned}$$

This grammar satisfies the conditions necessary for implementation as a DFA: it's right linear, and every production is of the form $a \rightarrow \epsilon$ or $b \rightarrow Xc$, where a , b , and c are nonterminals and X is a terminal. Consequently, you can represent this grammar with the state machine in Figure 4.1 and can parse that grammar using the state machine.

Figure 4.1. Using a State Machine for Parsing



State machines can't count.

The problem with a straightforward state-machine implementation of a real grammar is that a simple state machine can't count; or to be more precise, it can count only by adding additional states. Of course, most grammars can't count beyond 1 either. For example, the only way to get one, two, or three of something is with:

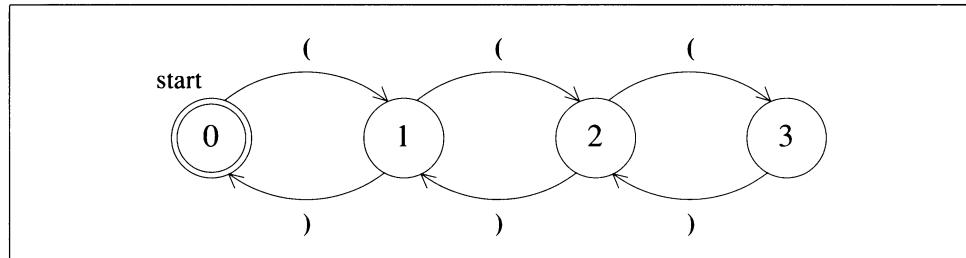
$$s \rightarrow \text{stooge} \mid \text{stooge stooge} \mid \text{stooge stooge stooge}$$

Nonetheless, a grammar *can* handle nested structures, such as parentheses, and keep track of the nesting level without difficulty, and a simple DFA cannot. For example, the following grammar recognizes properly nested lists of parentheses:

$$\begin{aligned}
 \text{list} &\rightarrow \text{plist } | \text{plist list} \\
 \text{plist} &\rightarrow (\text{list}) \mid \epsilon
 \end{aligned}$$

The grammar accepts input like the following: $(((())))$, but it rejects expressions without properly-nested parentheses. This grammar can't be parsed with a state machine alone, precisely because the state machine can't count. A straightforward state machine that recognizes nested parentheses is shown in Figure 4.2.

Figure 4.2. A State Machine to Recognize Nested Parentheses



Push-down automata (PDA).

This machine works fine as long as you don't have more than three levels of nesting. Each additional nesting level requires an additional state tacked on to the right edge of the machine, however. You don't want to limit the nesting level in the grammar itself, and it's not a great idea to modify the state machine on the fly by adding extra states as open parentheses are encountered. Fortunately, the situation can be rectified by using a state machine augmented with a stack. This state-machine/stack combination is called a *push-down automaton* or *PDA*. Most table-driven parsers are push-down automata. They are driven by a state machine and use a stack to keep track of the progress of that state machine.

A push-down automaton that recognizes nested parentheses is shown in Figure 4.3. The associated state table is in Table 4.1. A PDA is different from a normal state machine in several ways. First of all, the stack is an integral part of the machine. The number at the top of stack is the current state number. Secondly, the contents of the state table are not next-state numbers; rather, they are actions to perform, given a current state (on the top of stack) and input symbol. Four types of actions are supported:

- accept* A sentence in the input grammar has been recognized.
- error* A syntax error has been detected in the input.
- push N* Push *N* onto the stack, effectively changing the current state to *N*.
- pop* Pop one item from the stack, changing the current state to whatever state number is uncovered by the pop. The input is advanced with each push or pop.

Figure 4.3. A Push-Down Automata to Recognize Nested Parentheses

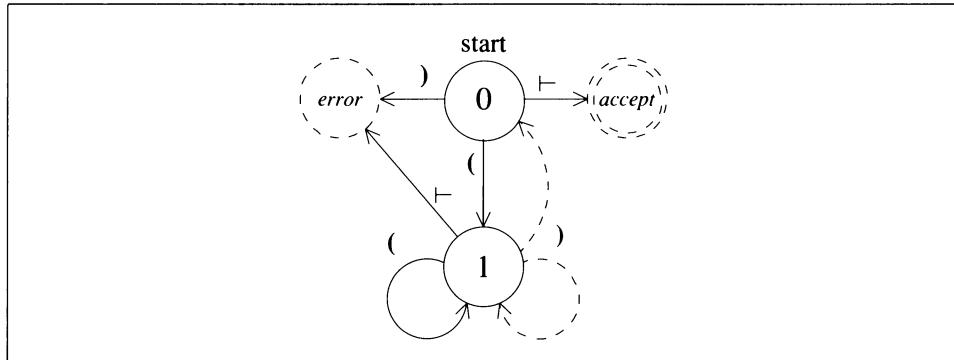


Table 4.1. State Table for PDA in Figure 4.3

		Input Symbol		
		()	⊤
State	0	push 1	error	accept
	1	push 1	pop	error

The following algorithm is used to parse an input sentence:

```

Push the number of the start state.
while( (action = state_table[ top_of_stack_symbol ][ input_symbol ]) ≠ accept )
{
    if( action = error )
        reject;
    else
        do the indicated action.
}
  
```

Using a PDA to count parentheses.

The parse stack and input stream for a parse of $((())()$) is summarized in Table 4.2. This PDA is using the stack as a counting device. The 0 marks the bottom of the stack, and the 1 is used as a nesting-level marker. The marker is pushed every time an open parenthesis is encountered and popped when the matching close parenthesis is found. The machine accepts an input string only if the start-state marker is on the stack when end of input is encountered.

Stack used for counting.

Table 4.2. A Parse of ((())())

Parse Stack	Input	Next Action
0	((())()) \vdash	push 1 and advance
0 1	(()()) \vdash	push 1 and advance
0 1 1) ()() \vdash	push 1 and advance
0 1	(()) \vdash	pop and advance
0 1 1	(()) \vdash	push 1 and advance
0 1 1 1) () \vdash	pop and advance
0 1 1) () \vdash	pop and advance
0 1) \vdash	pop and advance
0	\vdash	accept

This machine has several advantages over the previous one. Not only can it handle any level parenthesis nesting (or, at least, the only limitation in that regard is the depth of the parse stack), but it has many fewer states than would a straight DFA for a reasonable nesting level. We'll look at how to use pushdown automata in more sophisticated ways in a moment.

4.1.1 Recursive-Descent Parsers as Push-Down Automata*

How recursive-descent parsers use the stack.

It's instructive to see how the foregoing applies to a recursive-descent parser. Most programming languages use a run-time stack to control subroutine calls and returns. The simplest possible situation is an assembly-language, subroutine-call instruction like `JSR addr`, which pushes the address of the instruction following the `JSR` onto the run-time stack, and then starts executing instructions at `addr`. A matching `RET` instruction pops an address off the stack and starts executing code at that address. In this simple case, the stack frame consists of nothing but the return address.

A programming language that supports recursion must use the stack more extensively. The actual details are discussed in depth in Chapter Six, but for now, suffice it to say that space for subroutine arguments and automatic local variables are allocated from the stack as part of the subroutine-call process. A C subroutine call, for example, usually involves the following sequence of operations:

- (1) Push the arguments onto the stack.
- (2) Call the subroutine.
- (3) Allocate space for local variables by decrementing the stack pointer by some constant (I'm assuming that the stack grows towards low memory, as is usually the case).

Everything is done in reverse order when the subroutine returns, except that the arguments are discarded rather than being popped.

Stack frame.

This entire block (arguments, return address, and local variables) comprises the subroutine's *stack frame*. Figure 4.4 shows the various stack frames as they're created and destroyed as the naive recursive-descent parser in Chapter One processes the expression `1+2` using the grammar in Table 4.3. Each new stack frame represents a subroutine call, and when the stack frame disappears, the subroutine has returned. The input is shown on the right as it is gradually absorbed. In terms of compilers, the entire stack frame represents a nonterminal—the left-hand side of some production, and the local variables and argument components of the stack frame are the attributes associated with that nonterminal.

Table 4.3. An LL(1) Expression Grammar

1.	<i>stmt</i>	\rightarrow	ϵ
2.			<i>expr ; stmt</i>
3.	<i>expr</i>	\rightarrow	<i>term expr'</i>
4.	<i>expr'</i>	\rightarrow	$+ \text{ term } \text{expr}'$
5.			ϵ
6.	<i>term</i>	\rightarrow	<i>factor term'</i>
7.	<i>term'</i>	\rightarrow	$* \text{ factor } \text{term}'$
8.			ϵ
9.	<i>factor</i>	\rightarrow	number_or_id
10.			(<i>expr</i>)

Figure 4.4. The Parser's Stack Frames

statements							<i>I + 2 ;</i>
statements	expression						<i>I + 2 ;</i>
statements	expression	term					<i>I + 2 ;</i>
statements	expression	term	factor				<i>I + 2 ;</i>
statements	expression	term					<i>+ 2 ;</i>
statements	expression						<i>+ 2 ;</i>
statements	expression						<i>+ 2 ;</i>
statements	expression	expr_prime					<i>2 ;</i>
statements	expression	expr_prime	term				<i>2 ;</i>
statements	expression	expr_prime	term	factor			<i>;</i>
statements	expression	expr_prime	term				<i>;</i>
statements	expression	expr_prime					<i>;</i>
statements	expression	expr_prime	expr_prime				<i>;</i>
statements	expression	expr_prime	expr_prime				<i>;</i>
statements	expression						<i>;</i>
statements							

The question to ask is: what are these stack frames really doing as the parse progresses? They're doing two things. First, they're keeping track of the current position on the parse tree. The stack always holds the path from the root down to the current node on the tree. You can traverse up to a parent by popping one stack element. Second, the stack keeps track of attributes; since a stack frame represents a subroutine call, it also represents a nonterminal symbol in the parse tree. If you view a subroutine's arguments and local variables as attributes of the nonterminal, then the portion of the stack frame that holds these objects represents the attributes. The same nonterminal symbol can be on the stack in several places, but since each instance of the associated subroutine has its own stack frame (with its own local variables and arguments), each instance of the

Stack used to remember attributes and position on parse tree.

Recursive subroutines have unique stack frames, attributes.

Replace recursion with loop and explicit stack.

associated nonterminal can have a unique set of attributes attached to it. For example, when `expr_prime()` calls itself recursively (in the fifth stack from the bottom of Figure 4.4), there are two stack frames, each with a unique set of local variables and arguments.

You can make the transition from recursive descent to a table-driven parser by realizing that neither the recursion nor the subroutines themselves are actually necessary, provided that you can simulate the recursive-descent parser's stack frames using a local stack. That is, instead of using the implicit, run-time stack to keep track of the attributes and the position in the parse, you can maintain an explicit stack that does the same thing. That explicit stack can then be controlled by a nonrecursive subroutine and a table. For example, a production like this:

$$\begin{array}{lcl} \text{factor} & \rightarrow & \text{MINUS unop} \\ \text{unop} & \rightarrow & \text{NUMBER} \\ & | & \text{IDENTIFIER} \end{array}$$

can be implemented in a recursive-descent parser as follows:

```
factor()
{
    if( match( MINUS ) ) { advance(); unop(); }
    else                  { error(); }
}

unop()
{
    if      ( match( NUMBER      ) ) advance();
    else if( match( IDENTIFIER ) ) advance();
    else                  error();
}
```

And these subroutines can be simulated as follows:

```
#define factor 1
#define unop    2

push( factor );                      /* Push the goal symbol */
while( stack_not_empty() )
{
    switch( top_of_stack_symbol )
    {
        case factor: if( match(MINUS) ) { advance(); push( unop ); }
                      else                  { error(); }
                      pop();
                      break;

        case unop:   if      ( match( NUMBER      ) ) advance();
                      else if( match( IDENTIFIER ) ) advance();
                      else                  error();
                      pop();
                      break;
    }
}
```

Rather than calling a subroutine, a symbol representing that subroutine is pushed onto a local stack. The pops replace the `return` statements. Since the foregoing loop is, in effect, a DFA, you can translate it to a table with little difficulty.

4.2 Using a PDA for a Top-Down Parse*

A top-down parse can be implemented using a grammar, a stack, and the following algorithm. The parse stack is initialized by pushing the goal symbol (the topmost production in the grammar).

Rules for top-down parsing with a PDA.

- (0) If the parse stack is empty, the parse is complete.
- (1) If the item at top of stack is a nonterminal, replace it with its right-hand side, pushing the symbols in reverse order. For example, if you have a production: $a \rightarrow b c d$, and a is at the top of stack, pop the a and push the right-hand side in reverse order: first the d , then the c , and then the b . In the case of an ϵ production (a production with ϵ as its right-hand side), an item is popped, but nothing is pushed in its place. Goto (0).
- (2) Otherwise, if the item at top of stack is a terminal symbol, that symbol must also be the current lookahead symbol. If it's not, there's a syntax error, otherwise pop the terminal symbol and advance the input. Goto (0).

Note that step (2), above, requires certain symbols to be in the input at specific times because tokens at the top of stack must match the current input symbol. The parser is effectively predicting what the next input token will be. For this reason, the parser just described is often called a *predictive* parser. Note that only certain classes of grammars (which I'll discuss shortly) can be parsed by predictive parsers. Table 4.4 shows the parse stack as 1+2; is processed using the grammar in Table 4.3 on page 199 and a predictive parser.

Predictive parsers.

4.3 Error Recovery in a Top-Down Parser*

One of the advantages of top-down parsers is that effective error recovery is easy to implement. The basic strategy makes use of a set of tokens called a *synchronization set*. Symbols in this set are called *synchronization tokens*. Members of the synchronization set are typically symbols that can end blocks of code.

Synchronization set, tokens.

A syntax error occurs in a predictive parser when a token is at top of stack and that same token is not the current lookahead character. The synchronization set is used to recover from the error as follows:

- (1) Pop items off the parse stack until a member of the synchronization set is at top of stack. Error recovery is not possible if no such item is on the stack.
- (2) Read input symbols until the current lookahead symbol matches the symbol at top of stack or you reach end of input.
- (3) If you are at end of input, error recovery failed, otherwise you have recovered.

An alternate method, which is harder to implement but more robust, examines the parse stack and creates a list of those synchronization symbols that are on the stack when the error occurred. It then reads the input until any of these symbols are encountered, and pops items off the stack until the same symbol is at top of stack.

Synchronization-set stack.

Ideally, every production should have its own set of synchronization symbols, which are changed when the production is activated. To do this properly, you must have a stack of synchronization sets, and replacing a nonterminal with its right-hand side involves the following operations:

- (1) Push the associated synchronization set onto the synchronization-set stack.
- (2) Pop the nonterminal from the parse stack.
- (3) Push a special *pop* marker onto the parse stack.
- (4) Push the right-hand side of the nonterminal onto the parse stack in reverse order.

Table 4.4. A Top-Down Parse of $1+2;$

Parse Stack	Input	Comments
<i>stmt</i>	$1+2; \vdash$	Apply $\text{stmt} \rightarrow \text{expr} ; \text{stmt}$
—	$1+2; \vdash$	
<i>stmt</i>	$1+2; \vdash$	(This <i>stmt</i> is the one on the right of $\text{stmt} \rightarrow \text{expr} ; \text{stmt}$)
<i>stmt</i> ;	$1+2; \vdash$	
<i>stmt</i> ; <i>expr</i>	$1+2; \vdash$	Apply $\text{expr} \rightarrow \text{term expr}'$
<i>stmt</i> ;	$1+2; \vdash$	
<i>stmt</i> ; <i>expr'</i>	$1+2; \vdash$	
<i>stmt</i> ; <i>expr' term</i>	$1+2; \vdash$	Apply $\text{term} \rightarrow \text{factor term}'$
<i>stmt</i> ; <i>expr'</i>	$1+2; \vdash$	
<i>stmt</i> ; <i>expr' term'</i>	$1+2; \vdash$	
<i>stmt</i> ; <i>expr' term' factor</i>	$1+2; \vdash$	Apply $\text{factor} \rightarrow \text{num_or_id}$
<i>stmt</i> ; <i>expr' term'</i>	$1+2; \vdash$	
<i>stmt</i> ; <i>expr' term' num_or_id</i>	$1+2; \vdash$	TOS symbol matches lookahead, pop and advance
<i>stmt</i> ; <i>expr' term'</i>	$+2; \vdash$	Apply $\text{term}' \rightarrow \epsilon$
<i>stmt</i> ; <i>expr'</i>	$+2; \vdash$	Apply $\text{expr}' \rightarrow + \text{term expr}'$
<i>stmt</i> ;	$+2; \vdash$	
<i>stmt</i> ; <i>expr'</i>	$+2; \vdash$	
<i>stmt</i> ; <i>expr' term</i>	$+2; \vdash$	
<i>stmt</i> ; <i>expr' term +</i>	$+2; \vdash$	TOS symbol matches lookahead, pop and advance
<i>stmt</i> ; <i>expr' term</i>	$2; \vdash$	Apply $\text{term} \rightarrow \text{factor term}'$
<i>stmt</i> ; <i>expr'</i>	$2; \vdash$	
<i>stmt</i> ; <i>expr' term'</i>	$2; \vdash$	
<i>stmt</i> ; <i>expr' term' factor</i>	$2; \vdash$	Apply $\text{factor} \rightarrow \text{num_or_id}$
<i>stmt</i> ; <i>expr' term'</i>	$2; \vdash$	
<i>stmt</i> ; <i>expr' term' num_or_id</i>	$2; \vdash$	TOS symbol matches lookahead, pop and advance
<i>stmt</i> ; <i>expr' term'</i>	$; \vdash$	Apply $\text{term}' \rightarrow \epsilon$
<i>stmt</i> ; <i>expr'</i>	$; \vdash$	Apply $\text{expr}' \rightarrow \epsilon$
<i>stmt</i> ;	$; \vdash$	TOS symbol matches lookahead, pop and advance
<i>stmt</i>	$; \vdash$	Apply $\text{stmt} \rightarrow \epsilon$
—	$; \vdash$	Done

When a *pop* marker is found on the parse stack, one set is popped from the synchronization-set stack. The set at the top of the synchronization-set stack is used when an error is encountered.

In practice, a synchronization-set stack is usually not worth the effort to implement because most synchronization sets have the same tokens in them. It's usually adequate to have one synchronization set that can work in most situations and is used universally whenever an error is encountered. Taking C as an example, a good choice of synchronization tokens would be a semicolon, comma, close-brace, and close parenthesis. In other words, symbols that end commonly occurring phrases in the input language are good choices.

4.4 Augmented Grammars and Table-Driven Parsers*

Though the parser just described is nice, it's not very useful because there's no provision for code generation. It's like the initial stab at a recursive-descent parser back in Chapter One, before the code generation had been added to the actions. Fortunately this omission is easy to rectify.

As was discussed in Chapter Three, an *augmented* grammar is one in which code-generation actions are inserted directly into the grammar. An augmented version of the previous grammar is shown in Table 4.5. (It's the augmented grammar from the last chapter.) The augmented grammar is parsed just like the unaugmented one. The actions are just pushed onto the parse stack along with other symbols. When an action is at the top of the stack, it's executed.

Top-down code generation using augmented grammar.

Actions executed as parse tree traversed.

Table 4.5. An Augmented LL(1) Expression Grammar

1.	<i>statements</i>	\rightarrow	\vdash
2.			<i>expression ; statements</i>
3.	<i>expression</i>	\rightarrow	<i>term expression'</i>
4.	<i>expression'</i>	\rightarrow	$+ \text{term} \{ \text{op}('+' \text{)}; \} \text{expression}'$
5.			ϵ
6.	<i>term</i>	\rightarrow	<i>factor term'</i>
7.	<i>term'</i>	\rightarrow	$* \text{factor} \{ \text{op}('*' \text{)}; \} \text{term}'$
8.			ϵ
9.	<i>factor</i>	\rightarrow	number_or_id {create_tmp(yytext);}
10.			(<i>expression</i>)

The action subroutines that do the code generation (`op()` and `create_tmp()`) were discussed in Chapter Three, but to summarize:

- The `create_tmp()` subroutine generates code to copy the operand into an anonymous temporary variable, and it pushes the name onto a local Temporaries stack.
- The `op()` subroutine performs the operation on the two previously-generated temporary variables, popping their names off the stack, generating code to perform the operation, and then pushing the name of the temporary that holds the result.

A sample parse of $1+2$ is shown in Table 4.6. Note that the `op()` call follows the *term* in $\text{expr}' \rightarrow + \text{term} \{ \text{op}('+' \text{)}; \} \text{expr}'$. Since the temporaries are generated indirectly by the *terms*, two names will exist on the Temporaries stack when the `op()` call is executed.

4.4.1 Implementing Attributed Grammars in a PDA*

The previous example generated code without benefit of attributes. It is possible to use attributes in a PDA, however. A top-down parser such as the one we're looking at uses inherited attributes (the equivalent of subroutine arguments) exclusively. It's very difficult to represent return values (synthesized attributes) in this kind of parser, though the situation is reversed in the bottom-up parsers discussed in the next chapter.

Inherited attributes in top-down parser.

We'll support attributes by introducing a second stack, called an *attribute* or *value* stack, which simulates that part of a subroutine's stack frame that holds the subroutine arguments. This stack is a stack of structures, one field of which is used to hold the attribute associated with a production's left-hand side, and the other element of which is used to hold a symbol's own attribute. Every time a symbol is pushed onto the normal parse stack as part of a replacement operation, the attribute attached to the left-hand side that's being replaced is pushed onto the value stack (it's copied to both fields of the structure).

Value, attribute stack.

An action is executed when it's at the top of the stack, and all symbols to the right of the action in the production are still on the parse stack when the execution occurs. For example, when the `action()` in the following production is executed:

Table 4.6. A Top-Down Parse of $1+2$; Using the Augmented Grammar

Parse Stack	Input	Comments
<i>stmt</i>	1+2;↑	Apply $\text{stmt} \rightarrow \text{expr} ; \text{stmt}$
—	1+2;↑	
<i>stmt</i>	1+2;↑	
<i>stmt</i> ;	1+2;↑	
<i>stmt</i> ; <i>expr</i>	1+2;↑	Apply $\text{expr} \rightarrow \text{term expr}'$
<i>stmt</i> ;	1+2;↑	
<i>stmt</i> ; <i>expr'</i>	1+2;↑	
<i>stmt</i> ; <i>expr'</i> <i>term</i>	1+2;↑	Apply $\text{term} \rightarrow \text{factor term}'$
<i>stmt</i> ; <i>expr'</i>	1+2;↑	
<i>stmt</i> ; <i>expr'</i> <i>term'</i>	1+2;↑	
<i>stmt</i> ; <i>expr'</i> <i>term'</i> <i>factor</i>	1+2;↑	Apply $\text{factor} \rightarrow \mathbf{n} \{\text{create_tmp}(yytext);\}$
<i>stmt</i> ; <i>expr'</i> <i>term'</i>	1+2;↑	
<i>stmt</i> ; <i>expr'</i> <i>term'</i> { <i>create_tmp</i> (<i>yytext</i>);}	1+2;↑	
<i>stmt</i> ; <i>expr'</i> <i>term'</i> { <i>create_tmp</i> (<i>yytext</i>);} n	1+2;↑	TOS symbol matches lookahead, pop and advance
<i>stmt</i> ; <i>expr'</i> <i>term'</i> { <i>create_tmp</i> (<i>yytext</i>);}	+2;↑	Generate $t_0=1$;
<i>stmt</i> ; <i>expr'</i> <i>term'</i>	+2;↑	Apply $\text{term}' \rightarrow \epsilon$
<i>stmt</i> ; <i>expr'</i> <i>term'</i>	+2;↑	Apply $\text{expr}' \rightarrow + \text{term} \{\text{op}('+') ;\} \text{expr}'$
<i>stmt</i> ; <i>expr'</i>	+2;↑	
<i>stmt</i> ;	+2;↑	
<i>stmt</i> ; <i>expr'</i>	+2;↑	
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;}	+2;↑	
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term</i>	+2;↑	
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term</i> +	+2;↑	TOS symbol matches lookahead, pop and advance
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term</i>	+2;↑	Apply $\text{term} \rightarrow \text{factor term}'$
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;}	+2;↑	
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term'</i>	+2;↑	
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term'</i> <i>factor</i>	+2;↑	Apply $\text{factor} \rightarrow \mathbf{n} \{\text{create_tmp}(yytext);\}$
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term'</i>	+2;↑	
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term'</i> { <i>create_tmp</i> (<i>yytext</i>);}	+2;↑	
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term'</i> { <i>create_tmp</i> (<i>yytext</i>);} n	+2;↑	TOS symbol matches lookahead, pop and advance
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term'</i> { <i>create_tmp</i> (<i>yytext</i>);}	;↑	Generate $t_1=2$;
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;} <i>term'</i>	;↑	Apply $\text{term}' \rightarrow \epsilon$
<i>stmt</i> ; <i>expr'</i> { <i>op</i> ('+') ;}	;↑	Generate t_0+t_1 ;
<i>stmt</i> ; <i>expr'</i>	;↑	Apply $\text{expr}' \rightarrow \epsilon$
<i>stmt</i> ;	;↑	TOS symbol matches lookahead, pop and advance
<i>stmt</i>	↑	Apply $\text{stmt} \rightarrow \epsilon$
—	↑	Done

num_or_id is abbreviated as **n** so that the table can fit onto the page.

$$a \rightarrow b \{\text{action}();\} c d$$

The action symbol itself is at the top of the stack, *c* is directly underneath it, and *d* is under the *c*. The attribute for a symbol can be modified from within an action by modifying the right-hand-side field of an attribute-stack structure. The distance from the top of stack to the required structure is the same as the distance in the grammar from action to the required symbol. In the earlier example, *d* is two symbols to the right of the `{action()}` and its attributes are two symbols beneath the `{action()}`'s attributes on the value stack. If that action wants to modify the attribute for *d* it need only modify the attribute at offset 2 from the top of the attribute stack.¹ The value stack can be implemented using the modified top-down parse algorithm in Table 4.7.

1. It's also possible to keep track of attributes on the parse stack itself rather than on an auxiliary value stack. See [Lewis], pp. 311–337. This method is difficult to implement so is not discussed here.

Table 4.7. Top-Down Parse Algorithm with Attributes

<p>Data Structures: A <i>parse stack</i> of ints. A <i>value stack</i> of the following structures:</p> <pre>typedef struct { YYSTYPE left; YYSTYPE right; } yyvstype;</pre> <p>YYSTYPE is an arbitrary type—it is a character pointer in the current example. A variable, <i>lhs</i>, of type YYSTYPE.</p>
<p>Initially: Push the start symbol on the parse stack. Push garbage onto the value stack.</p>
<ul style="list-style-type: none"> (0) If the parse stack is empty, the parse is complete. (1) If the item at the top of the parse stack is an action symbol: execute the associated code and pop the item. (2) If the item at the top of the parse stack is a nonterminal: <ul style="list-style-type: none"> (a) <i>lhs</i> = the right field of the structure currently at the top of the value stack. (b) replace the symbol at the top of the parse stack with its right-hand side, pushing the symbols in reverse order. Every time an item is pushed onto the parse stack, push an item onto the value stack, initializing both fields to <i>lhs</i>. (c) Goto (0). (3) Otherwise, if the item at the top of the parse stack is a terminal symbol, that symbol must also be the current lookahead symbol. If it's not, there's a syntax error, otherwise pop the terminal symbol and advance the input. Goto (0).

I'll give an example of how to use this attribute-passing mechanism in a moment.

First, however, I need to introduce a new notation for attribute management. Attributes need only be referenced explicitly inside an action. Rather than giving attributes arbitrary names, all you need is a mechanism to reference them from the perspective of the action itself. The following notation, similar to the one used by **yacc**, does the job: The attribute associated with a left-hand side is always called **\$\$**, and attributes for symbols on the right-hand side are referenced by the notation **\$1**, **\$2**, and so forth, where the number represents the offset of the desired symbol to the right of the current action. Consider this production:

```
expr' → + { $1=$2=newname () ; }
term   { printf ("%s+=%s\n", $$, $0); freename ($0); }
expr'
```

(That's one long production, split up into several lines to make it easier to read, not three right-hand sides.) The **\$1** in **{ \$1=\$2=newname () ; }** attaches an attribute to the *term* because *term* is one symbol to the right of the action. When the parser replaces this *term* with its right-hand side, the symbols on that right-hand side can all use **\$\$** to access the

Attribute notation (**\$N**).

attribute associated with the *term*. In the current right-hand side, the $\$\$$ in the `printf()` statement references an attribute that was attached to the *expr'* at some earlier point in the parse. The $\$2$ in `{ $1=$2=newname(); }` references the second action, which is two symbols to the right of `{ $1=$2=newname(); }`. The second action can get at this attribute using $\$0$. That is, the value assigned to $\$2$ in the first action is accessed with $\$0$ in the second action.

LLama can translate the dollar notation directly into value-stack references when it outputs the `switch` statement that holds the action code. When an action is executed, $\$\$$ is the left-hand-side field of the structure at the top of the value stack, $\$0$ is the right-hand-side field of the same structure, $\$1$ is the structure at offset one from the top of the stack, $\$2$ is at offset two, and so forth.

The entire grammar, modified to use our new notation, is shown in Table 4.8. You can see how all this works by looking at the parse of `1+2` that is shown in Table 4.9. Remember that the symbols on the right-hand side of a replaced nonterminal inherit the attributes in the nonterminal's right field—both the `left` and `right` fields of the attribute-stack structure are initialized to hold the `right` field of the parent nonterminal.

Table 4.8. An Augmented, Attributed Grammar Using the \$ Notation

1.	<i>stmt</i>	$\rightarrow \epsilon$
2.		$ \quad \{ \$1=\$2=\text{newname}(); \} \text{expr} \{ \text{freename}(\$0); \} ; \text{stmt}$
3.	<i>expr</i>	$\rightarrow \text{term expr'}$
4.	<i>expr'</i>	$\rightarrow + \{ \$1=\$2=\text{newname}(); \} \text{term} \{ \text{printf}(" \%s+= \%s \n", \$\$, \$0); \text{freename}(\$0); \} \text{expr'}$
5.		$ \quad \epsilon$
6.	<i>term</i>	$\rightarrow \text{factor term'}$
7.	<i>term'</i>	$\rightarrow * \{ \$1=\$2=\text{newname}(); \} \text{factor} \{ \text{printf}(" \%s*= \%s \n", \$\$, \$0); \text{freename}(\$0); \} \text{term'}$
8.		$ \quad \epsilon$
9.	<i>factor</i>	$\rightarrow \text{number_or_id} \{ \text{printf}(" \%s=%0.*s \n", \$\$, \text{yylen}, \text{yytext}); \}$
10.		$ \quad (\text{expr})$

Table 4.9. A Parse of $1+2$ Showing Value Stacks

4.5 Automating the Top-Down Parse Process*

Now that we have an algorithm, we need a way to automate the parse process. The solution is a push-down automata in which a state machine controls the activity on the stack. I'll start by looking at how a table-driven top-down parser works in a general way, and then I'll examine a **LLama** output file in depth by way of example. I'll discuss how **LLama** itself works, and how the tables are generated, in the next section.

4.5.1 Top-Down Parse Tables*

PDA's in top-down parsing, parse tables.

The tables for the push-down automata used by the parser are relatively straightforward. The basic strategy is to give a numeric value to all the symbols in the grammar, and then make tables of these numeric values. For example, the numeric values in Table 4.10 represent the symbols in the augmented grammar in Table 4.5 on page 203.

Table 4.10. Numeric and Symbolic Values For Symbols in Expression Grammar

Symbolic Value	Numeric Value	Notes
LP	5	Terminal Symbols
NUM_OR_ID	4	
PLUS	2	
RP	6	
SEMI	1	
TIMES	3	
EOI	0	(End of input marker, created by LLama)
<i>expr</i>	257	Nonterminal Symbols
<i>expr'</i>	259	
<i>factor</i>	260	
<i>stmt</i>	256	
<i>term</i>	258	
<i>term'</i>	261	
{op('+');}	512	Actions
{op('*');}	513	
{create_tmp(yytext);}	514	

Numeric values for symbols.

The various symbol types can be distinguished by their numeric value. The end-of-input marker is always 0, tokens are all in the range 1–6, nonterminals are in the range 256–261, and actions are in the range 512–514. The parse stack contains these numbers rather than the terminals, nonterminals, and actions that they represent. For example, every appearance of an *expr* in a stack picture corresponds to a 257 on the real stack. The earlier parse of 1+2 is shown in both symbolic and numeric form in Table 4.11. For clarity, I will continue to use the symbolic names rather than the numbers, but remember that the stack itself actually has the numeric, not the symbolic, values on it.

Executing actions.

The parser takes care of actions with a switch statement controlled by symbolic constants—like the one in Listing 4.1.² The **switch** is executed when a number representing an action is encountered at top of stack—that same number (actnum) is used as the argument to the **switch**, and there's a **case** for every action symbol.

The parse itself is accomplished using two tables, *Yy_pushtab* and *Yy_d*, shown in Figures 4.5 and 4.6. The algorithm is in Table 4.12. *Yy_pushtab* holds the right-hand

2. This **switch** statement could be made easier to read by using macros to convert the numbers to a more readable string, but this mapping would make both **LLama** itself and the output code more complex, so I decided not to do it. The symbol table is available from **LLama** if you need to examine the output code.

Table 4.11. A Parse of 1+2 Showing Both Symbols and Numeric Equivalents

Stack (Symbols)	Stack (Numbers)
<i>stmt</i>	256
—	—
<i>stmt</i>	256
<i>stmt ;</i>	256 1
<i>stmt ; expr</i>	256 1 257
<i>stmt ;</i>	256 1
<i>stmt ; expr'</i>	256 1 259
<i>stmt ; expr' term</i>	256 1 259 258
<i>stmt ; expr'</i>	256 1 259
<i>stmt ; expr' term'</i>	256 1 259 261
<i>stmt ; expr' term' factor</i>	256 1 259 261 260
<i>stmt ; expr' term'</i>	256 1 259 261
<i>stmt ; expr' term' {2}</i>	256 1 259 261 514
<i>stmt ; expr' term' {2} n</i>	256 1 259 261 514 4
<i>stmt ; expr' term' {2}</i>	256 1 259 261 514
<i>stmt ; expr' term'</i>	256 1 259 261
<i>stmt ; expr'</i>	256 1 259
<i>stmt ;</i>	256 1
<i>stmt ; expr'</i>	256 1 259
<i>stmt ; expr' {0}</i>	256 1 259 512
<i>stmt ; expr' {0} term</i>	256 1 259 512 258
<i>stmt ; expr' {0} term +</i>	256 1 259 512 258 2
<i>stmt ; expr' {0} term</i>	256 1 259 512 258
<i>stmt ; expr' {0}</i>	256 1 259 512
<i>stmt ; expr' {0} term'</i>	256 1 259 512 261
<i>stmt ; expr' {0} term' factor</i>	256 1 259 512 261 260
<i>stmt ; expr' {0} term'</i>	256 1 259 512 261
<i>stmt ; expr' {0} term' {2}</i>	256 1 259 512 261 514
<i>stmt ; expr' {0} term' {2} n</i>	256 1 259 512 261 514 4
<i>stmt ; expr' {0} term' {2}</i>	256 1 259 512 261 514
<i>stmt ; expr' {0} term'</i>	256 1 259 512 261
<i>stmt ; expr' {0}</i>	256 1 259 512
<i>stmt ; expr'</i>	256 1 259
<i>stmt ;</i>	256 1
<i>stmt</i>	256
—	—
{0} represents \$1=\$2=newname ();	
{2} represents printf ("%s=%s\n", \$\$,\$0);freename (\$0);	

Listing 4.1. Executing Actions from a **switch**

```

1  switch( actnum )
2  {
3  case 512: { op('+'); }           break;
4  case 513: { op('*'); }           break;
5  case 514: { create_tmp(yytext); } break;
6
7  default:   printf("INTERNAL ERROR: Illegal action number\n");
8      break;
9 }
```

sides of the productions in reverse order (because you have to push them in reverse order when you do a replace operation). The index into the `Yy_pushtab` array is the production number—an arbitrary, but unique, number assigned to each production. Here, the topmost production in the grammar is Production 0, the next one is Production 1, and so

forth.³ The right-hand sides are represented by zero-terminated arrays of the earlier symbolic values. Yy_d is used to determine the actual replacement to perform. It is indexed by potential lookahead symbols along one axis, and nonterminals along the other, and is used with the parse algorithm in Table 4.12.

Figure 4.5. LLama Parse Tables: $\text{Yy_pushtab}[]$

Yy_pushtab						
•	0					Yyp00
•	256, 1, 257, 0 <i>stmt SEMI expr</i>					Yyp01
•	259, 258, 0 <i>expr' term</i>					Yyp02
•	259, 512, 258, 2, 0 <i>expr' {op(ADD);} term PLUS</i>					Yyp03
•	0					Yyp04
•	261, 260, 0 <i>term' factor</i>					Yyp05
•	261, 513, 260, 3, 0 <i>term' {op(MULTIPLY)} factor TIMES</i>					Yyp06
•	0					Yyp07
•	514, 4, 0 (rvalue(yytext)) NUM_OR_ID					Yyp08
•	6, 257, 5, 0 RP expr LP					Yyp09

Figure 4.6. LLama Parse Tables: $\text{Yy_d}[]$

Yy_d	\vdash	SEMI	PLUS	TIMES	NUM_OR_ID	LP	RP
<i>stmt</i>	0	-1	-1	-1	1	1	-1
<i>expr</i>	-1	-1	-1	-1	2	2	-1
<i>term</i>	-1	-1	-1	-1	5	5	-1
<i>expr'</i>	-1	4	3	-1	-1	-1	4
<i>factor</i>	-1	-1	-1	-1	8	8	-1
<i>term'</i>	-1	7	7	6	-1	-1	7

At this point, I'd suggest running through the parse of $1+2$ again, using the algorithm and tables just discussed.

3. Production numbers for **LLama**-generated parsers can be found by looking at the *llout.sym* file generated by **LLama**'s *-D*, *-S*, or *-s* switch.

Table 4.12. Top-Down Parse Algorithm, Final Version

<i>parse stack</i>	a stack of ints
<i>value stack</i>	a stack of yyvstype structures, with left and right fields, as described earlier. This stack is the same size as the parse stack.
<i>Yy_d</i>	one of the two tables in Figure 4.6.
<i>Yy_pushtab</i>	the other of the two tables in Figure 4.6.
Initially:	Push the numeric value of the goal symbol onto the parse stack. Push a meaningless place marker onto the value stack so that it is in phase with the parse stack.
	<pre> while(the parse stack is not empty) { if(the symbol at top of stack represents an action) { Execute the associated code and pop the action item. } else if(the symbol at top of stack represents a terminal) { if(that symbol doesn't match the input symbol) Syntax error. else { Pop one item off the parse stack. Advance. } } else /* top-of-stack symbol is a nonterminal */ { what_to_do = Yy_d[top-of-stack symbol] [current lookahead symbol]; if(what_to_do == -1) Syntax error. else { val = right field of item at top of value stack Pop one item off both the parse and value stacks. Push all the items listed in Yy_{pushtab}[what_to_do] onto the parse stack. Push the same number of items onto the value stack, initializing both fields to val. } } } </pre>

4.6 LL(1) Grammars and Their Limitations*

The concept of an LL(1) grammar was introduced in Chapter Three. This class of grammar can be parsed by a parser that reads the input from left to right and does a left-most derivation of the parse tree. It needs at most one character of lookahead. There are LL(0) grammars, which require no lookahead, but they're too limited to be useful in a compiler—all that they can do is recognize finite strings. LL(0) grammars are a subset

LL(0) grammars.

of the right-linear grammars discussed in the last chapter—they can have no ϵ productions in them.

The main task of a program like **LLama** is to manufacture the parse tables from the tokenized input grammar. The basic parse table (called Yy_d in the previous section) is a two-dimensional array, the rows are indexed by nonterminal symbols and columns by terminal symbols. The array contains either an error indicator or the number of the production to apply when a specified nonterminal is at the top of the parse stack and a particular terminal is in the input stream. This parse table can only be created if the input grammar is an LL(1) grammar.

How a parser decides which productions to apply.

An LL(1) grammar is limited in many ways, which are best explained by looking at how the parser determines what productions to apply in a given situation, and by looking at grammars that the parser can't use because it can't make this determination. Consider the following production that recognizes simple expressions comprised of numbers, parentheses and minus signs.

$$\begin{array}{lll} 1: & \textit{expression} & \rightarrow \texttt{OPEN_PAREN } \textit{expression} \texttt{ CLOSE_PAREN} \\ & 2: & | \quad \texttt{NUMBER MINUS } \textit{expression} \end{array}$$

If an *expression* is on top of the parse stack, the parser must determine which of the two possible productions to apply. It makes the decision by looking at the current lookahead token. If this token is a **OPEN_PAREN**, the parser applies the top production; if it's a **NUMBER**, the bottom production is applied. If any other symbol is the lookahead symbol, then an error has occurred. Note that if you were to add a production of the form

$$\textit{expression} \rightarrow \texttt{NUMBER PLUS } \textit{expression}$$

to the foregoing grammar, you could no longer determine which of the two right-hand sides that start with a **NUMBER** to apply when an *expression* was on top of the stack. The grammar would not be LL(1) in this case.

The situation is complicated somewhat in the following grammar:

$$\begin{array}{lll} 1: & \textit{expression} & \rightarrow \texttt{OPEN_PAREN } \textit{expression} \texttt{ CLOSE_PAREN} \\ 2: & & | \quad \textit{term MINUS } \textit{expression} \\ 3: & \textit{term} & \rightarrow \texttt{NUMBER} \\ 4: & & | \quad \texttt{IDENTIFIER} \end{array}$$

Here also, the top production is applied when an *expression* is on top of the stack and an **OPEN_PAREN** is in the input. You need to look further than the actual production to determine whether you can apply Production 2, however. Since the second production starts with a nonterminal, you can't tell what to do by looking at Production 2 only. You can resolve the problem by tracing down the grammar, looking at the symbols that can occur at the beginning of the nonterminal. Since a *term* begins the second production, and a *term* can start with either a **NUMBER** or **IDENTIFIER**, the parser can apply the second production if an *expression* is on top the stack and the input symbol is a **NUMBER** or **IDENTIFIER**. Note that problems would arise if you added a production of the form

$$\textit{expression} \rightarrow \texttt{NUMBER}$$

in the previous grammar, because, if a **NUMBER** were the lookahead symbol, the parser couldn't determine whether to apply Production 2 or the new production. If this situation existed, the grammar would not be LL(1).

Let's complicate the situation further with this grammar, which describes a simple compound statement that can either contain another statement or be empty:

1:	<i>statement</i>	→	OPEN_CURLY <i>expression</i> CLOSE_CURLY
2:			<i>expression</i> SEMICOLON
3:	<i>expression</i>	→	OPEN_PAREN <i>expression</i> CLOSE_PAREN
4:			<i>term</i> MINUS <i>expression</i>
5:			ϵ
6:	<i>term</i>	→	NUMBER
7:			IDENTIFIER

Production 1 is applied if a *statement* is on top of the stack and the input symbol is an **OPEN_CURLY**. Similarly, Production 2 is applied when a *statement* is on top of the stack and the input symbol is an **OPEN_PAREN**, **NUMBER**, or **IDENTIFIER** (an **OPEN_PAREN** because an *expression* can start with an **OPEN_PAREN** by Production 3, a **NUMBER** or **IDENTIFIER** because an *expression* can start with a *term*, which can, in turn, start with a **NUMBER** or **IDENTIFIER**). The situation is complicated when an *expression* is on top of the stack, however. You can use the same rules as before to figure out whether to apply Productions 3 or 4, but what about the ϵ production (Production 5)? The situation is resolved by looking at the symbols that can follow an *expression* in the grammar. If *expression* goes to ϵ , it effectively disappears from the current derivation (from the parse tree)—it becomes transparent. So, if an *expression* is on top of the stack, apply Production 5 if the current lookahead symbol can follow an *expression* (if it is a **CLOSE_CURLY**, **CLOSE_PAREN**, or **SEMICOLON**). In this last situation, there would be serious problems if **CLOSE_CURLY** could also start an *expression*. The grammar would not be LL(1) were this the case.

4.7 Making the Parse Tables*

I'll now formalize the foregoing rules and show how to use these rules to make a top-down parse table such as the one in Figure 4.6 on page 210. This section discusses the theory; code that implements this theory is presented later on in the chapter.

4.7.1 FIRST Sets*

The set of terminal symbols that can appear at the far left of any parse tree derived from a particular nonterminal is that nonterminal's **FIRST** set. Informally, if you had a bunch of productions that all had *expression* on the left-hand side, **FIRST**(*expression*) (pronounced “first of *expression*”) would comprise all terminals that can start an expression. In the case of the C language, this set would include numbers, open parentheses, minus signs, and so forth, but it would not include a close parenthesis, because that symbol can't start an expression. Note that ϵ is considered to be a terminal symbol, so it can appear in a **FIRST** set. First sets can be formed using the rules in Table 4.13. I'll demonstrate how to find **FIRST** sets with an example, using the grammar in table 4.13.

The **FIRST** sets are put together in a multiple-pass process, starting out with the easy ones. Initially, add those nonterminals that are at the far left of a right-hand side:

FIRST (<i>stmt</i>)	=	{ }
FIRST (<i>expr</i>)	=	{ }
FIRST (<i>expr'</i>)	=	{PLUS}
FIRST (<i>term</i>)	=	{ }
FIRST (<i>term'</i>)	=	{TIMES}

Computing **FIRST** sets,
an example.

Table 4.13. Finding FIRST Sets

- (1) $\text{FIRST}(A)$, where A is a terminal symbol, is $\{A\}$. If A is ϵ , then ϵ is put into the FIRST set.
 (2) Given a production of the form

$$s \rightarrow A \alpha$$

where s is a nonterminal symbol, A is a terminal symbol, and α is a collection of zero or more terminals and nonterminals, A is a member of $\text{FIRST}(s)$.

- (3) Given a production of the form

$$s \rightarrow b \alpha$$

where s and b are single nonterminal symbols, and α is a collection of terminals and nonterminals, everything in $\text{FIRST}(b)$ is also in $\text{FIRST}(s)$.

This rule can be generalized. Given a production of the form:

$$s \rightarrow \alpha \mathcal{B} \beta$$

where s is a nonterminal symbol, α is a collection of zero or more nullable nonterminals,[†] \mathcal{B} is a single terminal or nonterminal symbol, and β is a collection of terminals and nonterminals, then $\text{FIRST}(s)$ includes the union of $\text{FIRST}(\mathcal{B})$ and $\text{FIRST}(\alpha)$. For example, if α consists of the three nullable nonterminals x , y , and z , then $\text{FIRST}(s)$ includes all the members of $\text{FIRST}(x)$, $\text{FIRST}(y)$, and $\text{FIRST}(z)$, along with everything in $\text{FIRST}(\mathcal{B})$.

[†] A nonterminal is nullable if it can go to ϵ by some derivation. ϵ is always a member of a nullable nonterminal's FIRST set.

Table 4.14. Yet Another Expression Grammar

1:	<i>stmt</i>	\rightarrow	<i>expr SEMI</i>
2:	<i>expr</i>	\rightarrow	<i>term expr'</i>
3:			ϵ
4:	<i>expr'</i>	\rightarrow	PLUS <i>term expr'</i>
5:			ϵ
6:	<i>term</i>	\rightarrow	<i>factor term'</i>
7:	<i>term'</i>	\rightarrow	TIMES <i>factor term'</i>
8:			ϵ
9:	<i>factor</i>	\rightarrow	LEFT_PAREN <i>expr RIGHT_PAREN</i>
10:			NUMBER

$$\text{FIRST}(\text{factor}) = \{\text{LEFT_PAREN}, \text{NUMBER}\}$$

Next, close the sets (perform a closure operation on the initial sets) using the foregoing rules. Everything in $\text{FIRST}(\text{factor})$ is also in $\text{FIRST}(\text{term})$ because factor is the left-most symbol on the right-hand side of term . Similarly, everything in $\text{FIRST}(\text{term})$ is also in $\text{FIRST}(\text{expr})$, and everything in $\text{FIRST}(\text{expr})$ is also in $\text{FIRST}(\text{stmt})$. Finally, expr is a nullable nonterminal at the left of stmt 's right-hand side, so I'll add **SEMI** to $\text{FIRST}(\text{stmt})$. Applying these relationships yields the following first sets:

$$\begin{aligned} \text{FIRST}(\text{stmt}) &= \{\text{LEFT_PAREN}, \text{NUMBER}, \text{SEMI}\} \\ \text{FIRST}(\text{expr}) &= \{\text{LEFT_PAREN}, \text{NUMBER}\} \end{aligned}$$

$\text{FIRST}(expr')$	=	{PLUS}
$\text{FIRST}(term)$	=	{LEFT_PAREN, NUMBER}
$\text{FIRST}(term')$	=	{TIMES}
$\text{FIRST}(factor)$	=	{LEFT_PAREN, NUMBER}

One final note: the FIRST notation is a little confusing because you see it used in three ways:

- FIRST(A) (A is a terminal) is A. Since ϵ is a terminal, then $\text{FIRST}(\epsilon) = \{\epsilon\}$.
- FIRST(x) (x is a nonterminal) is that nonterminal's FIRST set, described above.
- FIRST(α) (α is a collection of terminals and nonterminals) is the FIRST set computed using the procedure in Rule (3), above: FIRST(α) always includes the FIRST set of the leftmost symbol in α . If that symbol is nullable, then it is the union of the FIRST sets of the first two symbols, if both of these symbols are nullable, then it is the union of the first three symbols, and so forth. If all the symbols in α are nullable, then FIRST(α) includes ϵ .

A subroutine that computes FIRST sets is presented towards the end of the current chapter in Listing 4.27 on page 305.

4.7.2 FOLLOW Sets*

The other set of symbols that you need in order to make the parse tables are the FOLLOW sets. A terminal symbol is in a nonterminal's FOLLOW set if it can follow that nonterminal in some derivation. You can find a nonterminal's FOLLOW set with the rules in Table 4.15. To see how Rule 3 in Table 4.14 works, consider the following grammar:

1: $compound_stmt$	\rightarrow	OPEN_CURLY stmt_list CLOSE_CURLY
2: $stmt_list$	\rightarrow	stmt_list stmt
3: $stmt$	\rightarrow	expr SEMI

CLOSE_CURLY is in FOLLOW(stmt_list) because it follows stmt_list in Production 1. CLOSE_CURLY is also in FOLLOW(stmt) because of the second of the following derivations:

$compound_stmt$	\xrightarrow{L}	OPEN_CURLY stmt_list CLOSE_CURLY
	\xrightarrow{L}	OPEN_CURLY stmt_list stmt CLOSE_CURLY
	\xrightarrow{L}	OPEN_CURLY stmt_list expr SEMI CLOSE_CURLY

The stmt_list in the first sentential form was replaced by the right-hand side of Production 2 (stmt_list stmt), and in that derivation, a CLOSE_CURLY followed the stmt.

I'll demonstrate how to compute FOLLOW sets with the earlier grammar in Table 4.14 on page 214. In the initial pass, apply the first two rules for forming FOLLOW sets: SEMI and RIGHT_PAREN are added to FOLLOW(expr) because they actually follow it in Productions 1 and 9; PLUS is added to FOLLOW(term) because everything in FIRST(expr') must be in FOLLOW(term) by Production 4; TIMES is added to FOLLOW(factor) because everything in FIRST(term') must be in FOLLOW(factor) by Productions 6 and 7. The initial pass looks like this:

$\text{FOLLOW}(stmt)$	=	{ \vdash }
$\text{FOLLOW}(expr)$	=	{SEMI, RIGHT_PAREN}
$\text{FOLLOW}(expr')$	=	{}

Computing FOLLOW sets, an example.

Table 4.15. Finding FOLLOW Sets

(1) If s is the goal symbol, \vdash (the end-of-input marker) is in $\text{FOLLOW}(s)$;

(2) Given a production of the form:

$$s \rightarrow \dots a \mathcal{B} \dots$$

where a is a nonterminal and \mathcal{B} is either a terminal or nonterminal, $\text{FIRST}(\mathcal{B})$ is in $\text{FOLLOW}(a)$;

To generalize further, given a production of the form:

$$s \rightarrow \dots a \alpha \mathcal{B} \dots$$

where s and a are nonterminals, α is a collection of zero or more nullable nonterminals and \mathcal{B} is either a terminal or nonterminal. $\text{FOLLOW}(a)$ includes the union of $\text{FIRST}(\alpha)$ and $\text{FIRST}(\mathcal{B})$.

(3) Given a production of the form:

$$s \rightarrow \dots a$$

where a is the rightmost nonterminal on the right-hand side of a production, everything in $\text{FOLLOW}(s)$ is also in $\text{FOLLOW}(a)$. (I'll describe how this works in a moment.) To generalize further, given a production of the form:

$$s \rightarrow \dots a \alpha$$

where s and a are nonterminals, and α is a collection of zero or more nullable nonterminals, everything in $\text{FOLLOW}(s)$ is also in $\text{FOLLOW}(a)$.

$$\text{FOLLOW}(\text{term}) = \{\text{PLUS}\}$$

$$\text{FOLLOW}(\text{term}') = \{\}$$

$$\text{FOLLOW}(\text{factor}) = \{\text{TIMES}\}$$

Now close the FOLLOW sets by making several passes through them, applying Rule 3 repetitively until nothing more is added to any FOLLOW set. The following holds:

- Everything in $\text{FOLLOW}(\text{expr})$ is also in $\text{FOLLOW}(\text{expr}')$ by Production 2.
- Everything in $\text{FOLLOW}(\text{term})$ is also in $\text{FOLLOW}(\text{term}')$ by Production 7.
- Since expr' is nullable, everything in $\text{FOLLOW}(\text{expr}')$ is also in $\text{FOLLOW}(\text{term})$ by Production 4.
- Since term' is nullable, everything in $\text{FOLLOW}(\text{term}')$ is also in $\text{FOLLOW}(\text{factor})$ by Production 7.

The first closure pass applies these identities to the original FOLLOW sets, yielding the following sets:

$$\text{FOLLOW}(\text{stmt}) = \{\vdash\}$$

$$\text{FOLLOW}(\text{expr}) = \{\text{SEMI}, \text{RIGHT_PAREN}\}$$

$$\text{FOLLOW}(\text{expr}') = \{\text{SEMI}, \text{RIGHT_PAREN}\}$$

$$\text{FOLLOW}(\text{term}) = \{\text{PLUS}, \text{SEMI}, \text{RIGHT_PAREN}\}$$

$$\text{FOLLOW}(\text{term}') = \{\text{PLUS}\}$$

$$\text{FOLLOW}(\text{factor}) = \{\text{TIMES}, \text{PLUS}\}$$

Another pass, using the same identities, adds a few more elements:

$$\text{FOLLOW}(\text{stmt}) = \{\vdash\}$$

$$\text{FOLLOW}(\text{expr}) = \{\text{SEMI}, \text{RIGHT_PAREN}\}$$

$\text{FOLLOW}(expr')$	=	{SEMI, RIGHT_PAREN}
$\text{FOLLOW}(term)$	=	{PLUS, SEMI, RIGHT_PAREN}
$\text{FOLLOW}(term')$	=	{PLUS, SEMI, RIGHT_PAREN}
$\text{FOLLOW}(factor)$	=	{TIMES, PLUS, SEMI, RIGHT_PAREN}

A third pass adds nothing to the FOLLOW sets, so you're done. A subroutine that computes FOLLOW sets is presented towards the end of the current chapter in Listing 4.28 on page 307.

4.7.3 LL(1) Selection Sets*

To review a bit, an LL(1) parse table looks like this:

<i>input symbol</i>	
<i>nonterminal</i>	error marker or production number

The columns are indexed by input symbol, the rows by nonterminal symbol. The table holds either a marker that signifies a syntax error or the number of a production to apply, given the current top-of-stack and input symbols. Each production has a unique, but arbitrary, production number. Typically, the start production is Production 0, the next one is Production 1, and so forth.

The LL(1) *selection set* for a given production is the set of nonterminals for which there are legal entries in any given row of the parse table. For example, a grammar could have the following productions in it:

1. $terminal \rightarrow \text{PERKIN_ELMER } pk$
2. | $\text{ADM3 } adm$
3. | dec_term
4. $dec_term \rightarrow \text{VT_52}$
5. | VT_100

LL(1) selection set.

The parse table for this grammar looks like this:

	PERKIN_ELMER	ADM3	VT_52	VT_100
terminal	1	2	3	3
dec_term	error	error	4	5

The number 1 is in the PERKIN_ELMER column of the terminal row, because Production 1 is applied if a term is at top of stack and PERKIN_ELMER is the lookahead symbol, and so forth. The same relationships are indicated by these selection sets:

- SELECT(1) = { PERKIN_ELMER }
- SELECT(2) = { ADM3 }
- SELECT(3) = { VT_52, VT_100 }
- SELECT(4) = { VT_52 }
- SELECT(5) = { VT_100 }

SELECT(3) indicates that Production 3 is selected if the left-hand side of Production 3 is on the top of the stack, and the current lookahead symbol is VT_52 or VT_100. In general, if you have both a production N (which takes the form $s \rightarrow \alpha$) and a token T, then T is in SELECT(N) if production N should be applied when s is on the top of the parse stack and T is the current lookahead symbol. Note that the selection sets are attached to the individual productions, not to the nonterminals. For a grammar to be LL(1), all productions that share a left-hand side must have unique selection sets, otherwise the parser wouldn't know what to do in a given situation. [This is the real definition of LL(1).] The

Table 4.16. Finding LL(1) Selection Sets

- | |
|---|
| <p>(1) A production is <i>nullable</i> if the entire right-hand side can go to ϵ. This is the case, both when the right-hand side consists only of ϵ, and when all symbols on the right-hand side can go to ϵ by some derivation.</p> <p>(2) For nonnullable productions: Given a production of the form</p> $s \rightarrow \alpha \mathcal{B} \dots$ <p>where s is a nonterminal, α is a collection of one or more nullable nonterminals, and \mathcal{B} is either a terminal or a nonnullable nonterminal (one that can't go to ϵ) followed by any number of additional symbols: the LL(1) select set for that production is the union of FIRST(α) and FIRST(\mathcal{B}). That is, it's the union of the FIRST sets for every nonterminal in α plus FIRST(\mathcal{B}). If α doesn't exist (there are no nullable nonterminals to the left of \mathcal{B}), then SELECT(s)=FIRST(\mathcal{B}).</p> <p>(3) For nullable productions: Given a production of the form</p> $s \rightarrow \alpha$ <p>where s is a nonterminal and α is a collection of zero or more nullable nonterminals (it can be ϵ): the LL(1) select set for that production is the union of FIRST(α) and FOLLOW(s). In plain words: if a production is nullable, it can be transparent—it can disappear entirely in some derivation (be replaced by an empty string). Consequently, if the production is transparent, you have to look through it to the symbols that can follow it to determine whether it can be applied in a given situation.</p> |
|---|

LL(1) selection sets are formed using the rules in Table 4.16.

Translating SELECT sets
into an LL(1) parse table.

Note that ϵ is used to compute the selection sets because s is nullable if FIRST(s) contains ϵ . Nonetheless, ϵ is not itself a member of any selection set. The selection sets can be translated into a parse table with the algorithm in Table 4.17.

Table 4.17. Translating SELECT Sets into LL(1) Parse Tables.

```

Initialize the table to all error transitions;
for( each production,  $N$ , in the grammar )
{
    lhs = the left-hand side of production  $N$ ;
    for( every token in SELECT( $N$ ) )
        parse_table[ lhs ][ token ] =  $N$ ;
}

```

4.8 Modifying Grammars*

All top-down parsers, including recursive-descent parsers, must use LL(1) grammars, which are quite limited. The main problem is that LL(1) grammars can't be left recursive. Since you can't just dispense with left associativity, LL(1) grammars would not be very useful unless some mechanism existed to translate left-recursive grammars into left associative, LL(1) grammars. This section looks at various ways that you can manipulate grammars algebraically to make them LL(1); the techniques are useful with other classes of grammars as well.

The following discussion uses the generic-grammar notation discussed at the end of Chapter Three without always calling out what everything means. Just remember that Greek letters represent collections of terminals and nonterminals, letters like \mathcal{A} and \mathcal{B} represent a single terminal or nonterminal, italics are used for single nonterminals, and boldface is used for single terminals.

4.8.1 Unreachable Productions*

An *unreachable* nonterminal is one that can't possibly appear in a parse tree rooted at the goal symbol. That is, there's no derivation from the goal symbol in which the unreachable nonterminal can appear in the viable prefix. For example, in the following grammar:

1. $s \rightarrow a$
2. $a \rightarrow \text{TERM}$
3. $b \rightarrow \text{TERM_TWO}$

Unreachable nonterminal.

Production 3 is clearly unreachable because b appears on no right-hand side. The situation is not always so cut and dried. Productions 3 and 4 are both unreachable in the following grammar, even though b and c both appear on right-hand sides:

1. $s \rightarrow a$
2. $a \rightarrow \text{TERM}$
3. $b \rightarrow \text{TERM_TWO } c$
4. $c \rightarrow \text{TERM_TOO } b$

Many of the transformation techniques discussed below create unreachable nonterminals and, since productions with these nonterminals on their left-hand sides have no useful function in the grammar, they should be removed. You can use the algorithm in Table 4.18.

4.8.2 Left Factoring*

A production like the following has two problems: It is clearly not LL(1) because the left-hand sides of both productions start with an **IF** token so they can't possibly have disjoint selection sets. Also, it is ambiguous because the same nonterminal appears twice on the right-hand side:

$$\begin{aligned} \text{statement} &\rightarrow \text{IF test THEN statement ELSE statement} \\ &| \quad \text{IF test THEN statement} \end{aligned}$$

Note that this is also a pretty serious ambiguity because it controls the ways that **if** and **else** statements bind. Input like the following:

Ambiguity in **if/else**, binding problems.

```
if( expr ) then
    if( expr2 ) then
        statement();
    else
        statement();
```

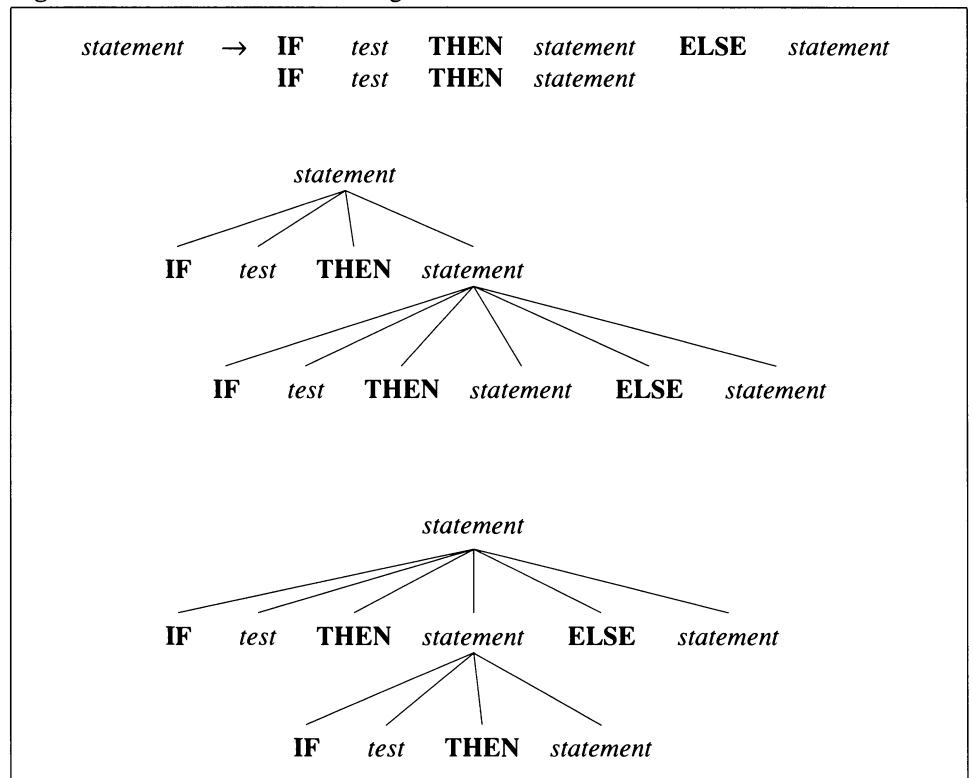
can create either of the trees pictured in Figure 4.7. The top parse causes the **else** to bind to the closest preceding **if**—the behavior required by most programming languages. In the second parser, though, the **else** incorrectly binds to first **if**.

Both problems can be solved by a process known as *left factoring*, which isolates the common parts of two productions into a single production. Any production of the form:

Left factoring.

Table 4.18. Eliminating Unreachable Productions

Data structures:	A stack A list of reachable nonterminals.
Initially:	Both the stack and the list are empty.
(1)	Add the goal symbol to the set of reachable nonterminals and push the goal symbol onto the stack.
(2)	while(the stack is not empty) { $s = \text{pop one item off the stack}$ for(each nonterminal, x , on a right-hand side of s) if(x is not in the list of reachable nonterminals) { push x ; add x to the list of reachable nonterminals } } }
(3)	Remove from the grammar all productions whose left-hand sides are not in the list of reachable nonterminals.

Figure 4.7. Parse Trees of an Ambiguous IF/ELSE Grammar

$$\begin{array}{lcl} a & \rightarrow & \alpha \beta_1 \\ & \dots & \\ & | & \alpha \beta_n \end{array}$$

where α is a sequence of one or more symbols that appear at the start of every right-hand side, and β_1 to β_n are dissimilar collections of zero or more symbols, can be replaced by the following:

$$\begin{array}{lcl} a & \rightarrow & \alpha a' \\ a' & \rightarrow & \beta_1 \\ & \dots & \\ & | & \beta_n \end{array}$$

The a' is just an arbitrary name for the new left-hand side. In the current example:

$$\begin{array}{ll} a & \text{corresponds to } statement \\ \alpha & \text{corresponds to } \mathbf{IF} \ test \ \mathbf{THEN} \ statement \\ \beta_1 & \text{corresponds to } \mathbf{ELSE} \ statement \\ \beta_2 & \text{corresponds to } \epsilon \end{array}$$

β_2 goes to ϵ because there's nothing in the second production that corresponds to the **ELSE statement** in the first production. You can plug the foregoing into our equation and replace the original productions with the following:

$$\begin{array}{ll} statement & \rightarrow \mathbf{IF} \ test \ \mathbf{THEN} \ statement \ opt_else_clause \\ opt_else_clause & \rightarrow \mathbf{ELSE} \ statement \\ & | \quad \epsilon \end{array}$$

Note that this replacement has also eliminated the ambiguity, because only one parse tree can now be generated from the earlier input.

4.8.3 Corner Substitution*

In general, given a left-hand side with one or more right-hand sides:

$$\begin{array}{lcl} p & \rightarrow & \alpha_1 \\ & | & \alpha_2 \\ & \dots & \\ & | & \alpha_n \end{array}$$

and given a production of the form:

$$s \rightarrow \beta \ p \ \gamma$$

this last production can be replaced with several productions of the form:

$$\begin{array}{lcl} s & \rightarrow & \beta \ \alpha_1 \ \gamma \\ & | & \beta \ \alpha_2 \ \gamma \\ & \dots & \\ & | & \beta \ \alpha_n \ \gamma \end{array}$$

This process is called a *substitution*.

Substitution.

In a production of the form $a \rightarrow \alpha$, the leftmost symbol on the right-hand side (α) is said to be a *corner* of a . An ϵ production doesn't have a corner. In a *corner substitution* you replace one or more of the nonterminal corners in a grammar with that corner's right-hand sides.

Corner.

For example, consider the following grammar, which recognizes a list of one or more ATOMs. The ATOMs can be single ATOMs, they can be followed by a bracketed number (LB and RB stand for left and right bracket), or they can be preceded by any

number of STARS.

1. $atom_list \rightarrow ATOM$
2. $| list_ele atom_list$
3. $list_ele \rightarrow ATOM\ LB\ NUMBER\ RB$
4. $| STAR\ list_ele$

This grammar is not LL(1) because, ATOM is in FIRST($list_ele$) by Production 3, and as a consequence, is also in SELECT(2). A substitution of the $list_ele$ corner of Production 2 can help fix this situation. First, you replace all instances of $list_ele$ in Production 2 with the right-hand sides of the two $list_ele$ productions (3 and 4), yielding:

1. $atom_list \rightarrow ATOM$
- 2a. $| ATOM\ LB\ NUMBER\ LB\ atom_list$
- 2b. $| STAR\ list_ele\ atom_list$
3. $list_ele \rightarrow ATOM\ LB\ NUMBER\ RB$
4. $| STAR\ list_ele$

The grammar is still not LL(1) because Productions 1 and 2a both start with an ATOM, but that situation can be rectified by left factoring:

1. $atom_list \rightarrow ATOM\ atom_list'$
- 2b. $| STAR\ list_ele\ atom_list$
- 1b. $atom_list' \rightarrow LB\ NUMBER\ RB\ atom_list$
- 1c. $| \epsilon$
3. $list_ele \rightarrow ATOM\ LB\ NUMBER\ RB$
4. $| STAR\ list_ele$

Corner substitution preserves LL(1) properties.

Q Grammars.

Corner substitution preserves the LL(1) properties of a grammar—grammars that start out LL(1) are still LL(1) after the corner substitution is made. You can see why if you consider how FIRST sets are computed. If x is a corner of production s , then $\text{FIRST}(s)$ includes $\text{FIRST}(x)$, and $\text{FIRST}(x)$ is computed by looking at the right-hand side of x . Substituting x for its right-hand side does not affect this computation.

One common use of corner substitution and left factoring takes advantage of this property by rearranging the grammar so that every right-hand side in the grammar either starts with a unique terminal symbol or is ϵ . That is, if a nonterminal has several right-hand sides, then all those right-hand sides will start with different terminal symbols or be ϵ . This particular type of grammar is called a Q grammar and is handy when you build recursive-descent parsers because it's very easy to code from a Q grammar. That is, given a Q grammar like the following:

$$\begin{array}{lcl} p & \rightarrow & T_1 \alpha_1 \\ & | & T_2 \alpha_2 \\ & \dots & \\ & | & T_n \alpha_n \\ & | & \epsilon \end{array}$$

where $T_1 \dots T_n$ are unique terminal symbols, you can code it like this:

```

p()
{
    switch( lookahead_character )
    {
        case T1: advance(); α1
        case T2: advance(); α2
        ...
        case Tn: advance(); αn

        default:      /* Handle the epsilon production */
    }
}

```

4.8.4 Singleton Substitution*

Generally, it's not productive to replace symbols that aren't corners—you shouldn't replace nonterminals that aren't at the far left of a right-hand side. The problem here is that the substitution usually creates several productions that have the same corner, and as a consequence the resulting grammar won't be LL(1). For example, substituting the *num_or_id* for its right-hand side in the following grammar:

1. $expr \rightarrow UNOP\ num_or_id$
2. $num_or_id \rightarrow NAME$
3. $num_or_id \rightarrow IDENTIFIER$

yields the following, non-LL(1) grammar:

1. $expr \rightarrow UNOP\ NAME$
- 1a. $expr \rightarrow UNOP\ IDENTIFIER$

You've just done a reverse left factoring.

If a production has only one right-hand side, of course, the substitution is harmless. In fact, it can reduce the size of the grammar by eliminating productions and is often useful for this reason. This sort of production is called a *singleton*.

It's sometimes desirable to create singletons by breaking out a group of symbols from the middle of a production. That is, the following:

$$s \rightarrow \alpha \beta \gamma$$

can be changed into:

$$\begin{aligned} s &\rightarrow \alpha s' \gamma \\ s' &\rightarrow \beta \end{aligned}$$

This technique is used primarily in bottom-up parsers, in which action symbols must be at the far right of the production. As a consequence, a production like this:

$$s \rightarrow tempest \{act(1);\} scene_5$$

must be implemented like this:

$$\begin{aligned} s &\rightarrow tempest\ s'\ scene_5 \\ s' &\rightarrow \{act(1);\} \end{aligned}$$

Substitutions can create ambiguity.

Singleton substitution used to isolate actions for bottom-up parser.

4.8.5 Eliminating Ambiguity*

Ambiguous productions are those that have more than one occurrence of a given nonterminal on their right-hand side. As we've seen, left factoring can be used to

eliminate ambiguity, moving the rightmost of these nonterminals further down the parse tree by introducing a new production. In general given an ambiguous production of the form:

$$s \rightarrow \alpha p \beta p \gamma$$

you can eliminate the ambiguity by introducing a new production, as follows:

$$\begin{array}{l} s \rightarrow \alpha p \beta s' \\ s' \rightarrow p \gamma \end{array}$$

If the grammar has a production of the form:

$$p \rightarrow s$$

Controlling associativity
when eliminating ambiguity.

this transformation makes the grammar left-associative, because there is now an indirect left recursion, demonstrated with the following derivation:

$$s' \Rightarrow p \gamma \Rightarrow s \gamma$$

If an ambiguous right-hand side is one of several, then all of these right-hand sides must move as part of the substitution. For example, given:

$$\begin{array}{l} e \rightarrow e + e \\ | \quad \text{NUM} \end{array}$$

you transform the grammar to be left-associative like this:

$$\begin{array}{l} e \rightarrow e + t \mid t \\ t \rightarrow e \\ | \quad \text{NUM} \end{array}$$

or right-associative like this:

$$\begin{array}{l} e \rightarrow t + e \mid t \\ t \rightarrow e \\ | \quad \text{NUM} \end{array}$$

Note that the foregoing transformation introduced a production of the form $e \rightarrow t$ into the grammar. You can use this production to modify the grammar again:

$$\begin{array}{l} e \rightarrow t + e \mid t \\ t \rightarrow t \\ | \quad \text{NUM} \end{array}$$

A production of the form $t \rightarrow t$ does nothing useful, so it can be eliminated:

$$\begin{array}{l} e \rightarrow t + e \mid t \\ t \rightarrow \text{NUM} \end{array}$$

To generalize, given grammars of the form:

$$s \rightarrow \alpha s \beta s \gamma \mid \alpha_1 \mid \dots \mid \alpha_n$$

you can eliminate the ambiguity in favor of right associativity by replacing the foregoing with the following:

$$\begin{array}{l} s \rightarrow \alpha s' \beta s \gamma \mid s' \\ s' \rightarrow s \mid \alpha_1 \mid \dots \mid \alpha_n \end{array}$$

Now, because a production of the form $s \rightarrow s'$ has been created you can substitute the s 's for s 's in the new production, yielding:

$$\begin{array}{lcl} s & \rightarrow & \alpha s' \beta s \gamma \mid s' \\ s' & \rightarrow & s' \mid \alpha_1 \mid \dots \mid \alpha_n \end{array}$$

and since that production of the form $s' \rightarrow s'$ doesn't do anything useful, it can be eliminated, yielding:

$$\begin{array}{lcl} s & \rightarrow & \alpha s' \beta s \gamma \mid s' \\ s' & \rightarrow & \alpha_1 \mid \dots \mid \alpha_n \end{array}$$

The same applies in a left-associative situation. Productions of the form:

$$s \rightarrow \alpha s \beta s \gamma \mid \alpha_1 \mid \dots \mid \alpha_n$$

can be disambiguated in favor of left-associativity by replacing the right occurrence of s rather than the left one in the initial step.

If a grammar has several ambiguous productions, it can be modified by repetitive application of the previous transformation. Similarly, if the s in the earlier general example appears on more than one right-hand side, then it can be replaced with s' in all the right-hand sides. For example, starting with the following expression grammar:

$$\begin{array}{lcl} 1. & e & \rightarrow e + e \\ 2. & & \mid e^* e \\ 3. & & \mid \text{NUM} \end{array}$$

I'll eliminate the ambiguity one production at a time. Note that the order in which the disambiguating rules are applied also affects operator precedence—the operators that are lower down in the grammar are of higher precedence. In an expression grammar like the current one, you want to start with the right-hand side that has the lowest-precedence operator, Production 1. Since addition must be left associative, the transformed production must also be left associative, so you must replace the right e in Production 1, yielding:

$$\begin{array}{lcl} 1a. & e & \rightarrow e + e' \\ 1b. & & \mid e' \\ 1c. & e' & \rightarrow e \\ 2. & & \mid e^* e \\ 3. & & \mid \text{NUM} \end{array}$$

Controlling precedence when eliminating ambiguity.

Now, since a production of the form $e \rightarrow e'$ exists, you can substitute e 's for e 's in Productions 1c, and 2, yielding:

$$\begin{array}{lcl} 1a. & e & \rightarrow e + e' \\ 1b. & & \mid e' \\ 1c. & e' & \rightarrow e' \\ 2. & & \mid e' * e' \\ 3. & & \mid \text{NUM} \end{array}$$

Production 1c now does nothing useful because the left- and right-hand sides are identical, so it can be removed from the grammar:

$$\begin{array}{lcl} 1a. & e & \rightarrow e + e' \\ 1b. & & \mid e' \\ 2. & e' & \rightarrow e' * e' \\ 3. & & \mid \text{NUM} \end{array}$$

You can now apply the same process to Production 2 to get the following grammar:

1a.	e	\rightarrow	$e + e'$
1b.			e'
2a.	e'	\rightarrow	$e' * e''$
2b.			e''
2c.	e''	\rightarrow	e'
3.			NUM

and substituting the e' for e'' in Production 2c, you get:

1a.	e	\rightarrow	$e + e'$
1b.			e'
2a.	e'	\rightarrow	$e' * e''$
2b.			e''
2c.	e''	\rightarrow	e''
3.			NUM

Finally, removing the redundant production yields:

1a.	e	\rightarrow	$e + e'$
1b.			e'
2a.	e'	\rightarrow	$e' * e''$
2b.			e''
3.	e''	\rightarrow	NUM

This grammar is left associative, and $+$ is lower precedence than $*$.

4.8.6 Eliminating Left Recursion*

LL(1) grammars cannot be left recursive.

LL(1) grammars cannot be left recursive. The basic proof of this statement is as follows:

- (1) In all practical grammars, if any right-hand side of a production is left recursive, there must be at least one nonrecursive right-hand side for the same nonterminal. For example, if a grammar has a singleton production of the form $s \rightarrow s\alpha$, all derivations that use that production go on forever:

$$\begin{aligned} s &\Rightarrow s\alpha \\ &\Rightarrow s\alpha\alpha \\ &\Rightarrow s\alpha\alpha\alpha \\ &\dots \end{aligned}$$

- You need to add a nonrecursive right-hand side of s for the derivation to terminate.
- (2) If x is a left-recursive nonterminal, the selection sets for all productions with x on their left-hand side must contain $\text{FIRST}(x)$, so, in all practical grammars, left-recursive productions must have selection-set conflicts.

Unfortunately, you need left recursion to get left associativity in a list. Fortunately, you can always eliminate left recursion from a grammar in such a way that the translated grammar recognizes the same input strings. The basic strategy looks at what left-recursive list grammars actually do. A list such as an expression (a list of operands separated by operators) is either a single operand or a single operand followed by one or more operator/operand pairs. You can express this structure with the following left-recursive grammar:

$$\begin{aligned} \text{list} &\rightarrow \text{operand} \\ &| \text{ list operator operand} \end{aligned}$$

The following nonrecursive grammar recognizes the same input strings, however:

$$\begin{array}{lcl} \text{list} & \rightarrow & \text{operand list}' \\ \text{list}' & \rightarrow & \text{operator operand} \\ | & & \epsilon \end{array}$$

So, our goal is to translate lists of the first form into lists of the second form. If a production is a self left recursive like the following:

$$s \rightarrow s \alpha | \beta$$

(the left-hand side is duplicated as the first symbol on the right-hand side: s is a single nonterminal; α and β are collections of terminals and nonterminals), you can make it nonrecursive by shuffling the production around as follows:

$$\begin{array}{lcl} s & \rightarrow & \beta s' \\ s' & \rightarrow & \alpha s' \\ | & & \epsilon \end{array}$$

Applying this translation rule to the following productions:

$$\text{expr} \rightarrow \text{expr} + \text{term}\{\text{act2}\} \mid \text{term}\{\text{act1}\}$$

the following relationships exist between the s , α , and β in the rule and the real productions:

$$\begin{array}{lcl} s & \rightarrow & s \quad \alpha \quad | \quad \beta \\ \text{expr} & \rightarrow & \text{expr} \quad + \text{term}\{\text{act2}\} \quad | \quad \text{term}\{\text{act1}\} \end{array}$$

s , which is a single nonterminal, corresponds to an *expression*; α , which is a collection of terminals and nonterminals, corresponds to the $+ \text{term}$, and β , which is also a collection of terminals and nonterminals, corresponds to the second *term* (in this case, it's a collection of only one symbol). Productions like the foregoing can be shuffled around to look like this:

$$\begin{array}{lcl} s & \rightarrow & \beta s' \\ s' & \rightarrow & \alpha s' \\ | & & \epsilon \end{array}$$

The s' is a new nonterminal. You could call it anything, but just adding a ' is easiest.

Applying the foregoing to the real productions yields:

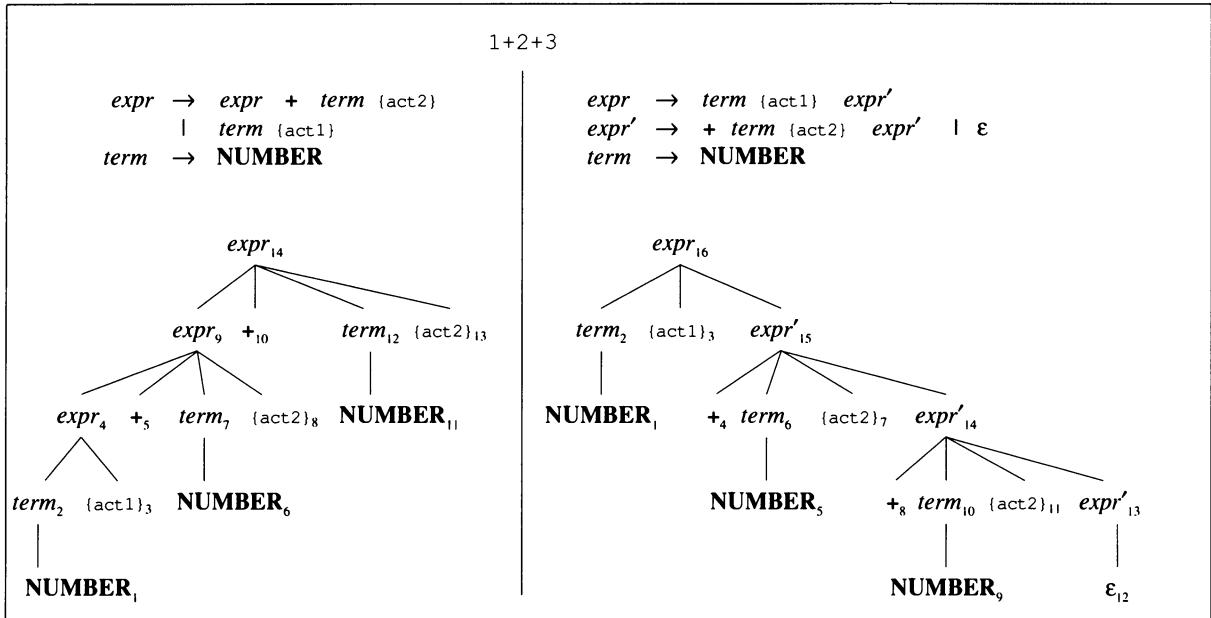
$$\begin{array}{llll} s & \rightarrow & \beta & s' \\ \text{expr} & \rightarrow & \text{term}\{\text{act1}\} & \text{expr}' \\ \\ s' & \rightarrow & \alpha & s' \quad | \quad \epsilon \\ \text{expr}' & \rightarrow & + \text{term}\{\text{act2}\} & \text{expr}' \quad | \quad \epsilon \end{array}$$

Figure 4.8 shows a parse for the input $1+2+3$ for both the untranslated and translated grammars. The important thing to notice is that, even though the parse tree for the translated grammar is now right associative (as you would expect from looking at the grammar), the order in which the terminal nodes are processed is identical in both trees. That is, the actions are performed in the same sequence relative to the input in both grammars. So there is still left associativity from the point of view of code generation, even though the parse itself is right associative.

Order in which actions
are executed with
transformed grammar.

The foregoing substitution process can be generalized for more complex grammars as follows:

Figure 4.8. Parse Trees for Translated and Untranslated, Left-Associative Grammars



$$s \rightarrow s \alpha_1 | s \alpha_2 | \dots | s \alpha_n | \beta_1 | \beta_2 | \dots | \beta_m$$

you can replace the foregoing with:

$$\begin{aligned} s &\rightarrow \beta_1 s' | \beta_2 s' | \dots | \beta_m s' \\ s' &\rightarrow \alpha_1 s' | \alpha_2 s' | \dots | \alpha_n s' | \epsilon \end{aligned}$$

Corner substitution translates indirect recursion to self recursion.

Though the method just described works only with self-recursive productions, it's an easy matter to use a corner substitution to make an indirectly recursive grammar self recursive. For example, in a grammar like this:

1. *expr* → *ele_list*
2. *ele_list* → **NUMBER**
3. | *expr PLUS NUMBER*

you can use a corner substitution to make the grammar self-recursive, replacing the *ele_list* corner in Production 1 with the right-hand sides of Productions 2 and 3 to yield the following:

- 1a. *expr* → **NUMBER**
- 1b. | *expr PLUS NUMBER*
2. *ele_list* → **NUMBER**
3. | *expr PLUS NUMBER*

Productions 2 and 3 are now unreachable, so can be eliminated from the grammar.

4.9 Implementing LL(1) Parsers

The remainder of this chapter discusses implementation details for the foregoing theory. Skip to the next chapter if you’re not interested.

4.9.1 Top-Down, Table-Driven Parsing—The LLama Output File

First, an in-depth look at the **LLama** output file, in which the top-down parse algorithm is actually implemented, seems in order. A user’s manual for **LLama** is presented in Appendix E, and you should read that appendix before continuing. This section examines the **LLama** output file created from the input file at the end of Appendix E. Listing 4.2 shows token definitions that are generated from the `%term` directives in the **LLama** input file. These are the same symbolic values that are used in the previous section.

Definitions generated from `%term`.

Listing 4.2. *llout.h*—Symbolic Values for Tokens

```

1 #define _EOI_      0
2 #define PLUS      1
3 #define TIMES     2
4 #define NUM_OR_ID 3
5 #define LP        4
6 #define RP        5
7 #define SEMI     6

```

The **LLama** C-source-code output file begins in Listing 4.3. Most of the file is copied from the template file *llama.par*, in a manner identical to the L^EX template file described in Chapter Three⁴. Lines three to 23 were copied from the input file, and the remainder of the listing was generated by **LLama**. Lines 28 to 33 define the boundaries of the numbers that represent terminals, nonterminals, and actions on the stack (as were described in the previous section). The minimums won’t change, but the maximums will.

Numerical limits of tokenized input symbols.

The next part of *llout.c* (in Listing 4.4) is copied from the template file. `<stdarg.h>` is included on Line 42 only if it hasn’t been included earlier. The `<tools/yystack.h>` file included on the next line contains various stack-manipulation macros, and is described in Appendix A. The macros on lines 45 to 47 determine what *x* is, using the earlier limits. The remainder of the listing is macros that define various parser limits described in Appendix E. These definitions are active only if you haven’t defined the macros yourself in a `%{ %}` block at the top of the input file. `YY_TTYPE` (on line 34) is used to declare the tables.

Parse- and value-stack declarations.

Various stacks are declared on lines 89 to 108 of Listing 4.4. The `<yystack.h>` macros described in Appendix A are used for this purpose, and various `<yystack.h>` macros are customized on lines 90 to 94 for the current use—`yyerror()` works like `printf()` but doesn’t mess up the windows when it’s used. The parse stack is defined on line 98, and the value to stack is on lines 100 to 108. Note that the `<yystack.h>` macros can’t be used here because the value stack is a stack of structures.

Parse-table declarations.

The next part of *llout.c* (in Listing 4.5) is generated by **LLama**. It contains the definitions for the parse tables discussed in the previous section. The tables in Listing

4. That’s why some of the following listings are labeled *llama.par* and others are labeled *llout.c*. The tables (which are generated by **LLama**) are in listings labeled *llout.c*. I’m sorry if this is confusing, but it’s the only way I can keep track of what’s where.

Listing 4.3. *llout.c*—File Header—Numeric Limits

```

1 #include <stdio.h>
2 #define YYDEBUG
3 /*-----
4  * Temporary-variable names are stored on a stack. name() pops a name off
5  * the stack and freename(x) puts it back. A real compiler would do some
6  * checking for stack overflow here, but there's no point in cluttering the
7  * code for now.
8 */
9
10 char *Namepool[] =
11 {
12     "t0", "t1", "t2", "t3", "t4", "t5", "t6", "t7", "t8", "t9"
13 };
14
15 char **Namep = Namepool ;
16
17 char *newname() { return (*Namep++); }
18 char *freename( char *x) { return (*--Namep = x); }
19
20 extern char *yytext;
21 extern int yyleng;
22
23 #define YYSTYPE char*
24
25 /*****/
26 #include "llout.h"
27
28 #define YY_MINTERM    1      /* Smallest terminal.          */
29 #define YY_MAXTERM    6      /* Largest terminal.          */
30 #define YY_MINNONTTERM 256   /* Smallest nonterminal.      */
31 #define YY_MAXNONTTERM 261   /* Largest nonterminal.      */
32 #define YY_START_STATE 256   /* Goal symbol (push to start parse). */
33 #define YY_MINACT     512   /* Smallest action.          */

```

Listing 4.4. *llama.par*—File Header—Macro Definitions

```

34 typedef unsigned char YY_TTYPE;           /* Type used for tables.      */
35 #define YYF      (YY_TTYPE)( -1 )        /* Failure transition in table. */
36
37 /*****
38  * Parser for llama-generated tables
39 */
40
41 #ifndef va_start
42 #include <stdarg.h>
43 #endif
44
45 #define YY_ISTERM(x)    (YY_MINTERM    <= (x) && (x) <= YY_MAXTERM    )
46 #define YY_ISNONTTERM(x) (YY_MINNONTTERM <= (x) && (x) <= YY_MAXNONTTERM)
47 #define YY_ISACT(x)      (YY_MINACT    <= (x)                      )
48
49 #ifndef YYACCEPT
50 #  define YYACCEPT return(0) /* Action taken when input is accepted. */
51 #endif
52

```

Listing 4.4. continued...

```

53 #ifndef YYABORT
54 # define YYABORT return(1) /* Action taken when input is rejected. */
55 #endif
56
57 #ifndef YYPRIVATE
58 # define YYPRIVATE static /* Define to a null string to make public. */
59 #endif
60
61 #ifndef YYMAXERR
62 # define YYMAXERR 25 /* Abort after this many errors. */
63 #endif
64
65 #ifndef YYMAXDEPTH
66 # define YYMAXDEPTH 128 /* Parse- and value-stack depth. */
67 #endif
68
69 #ifndef YYSTYPE
70 # define YYSTYPE int /* Used to declare fields in value stack. */
71 #endif
72
73 extern int yylineno; /* Supplied by LeX. */
74 extern int yyleng;
75 extern char *yytext;
76
77 extern int ii_plineno(); /* in l.lib::input.c. */
78 extern char *ii_ptext();
79 extern int ii_linenr();
80 extern char *ii_text();
81 extern void ii_mark_prev();
82
83 void yyerror( char *fmt, ... );
84
85 /-----
86 * Parse and value stacks:
87 */
88
89 #include <tools/yystack.h>
90 #undef yystk_cls
91 #define yystk_cls YYPRIVATE
92 #undef yystk_err
93 #define yystk_err(o) ((o) ? (yyerror("Stack overflow\n"),exit(1)) \
94 : (yyerror("Stack underflow\n"),exit(1)))
95
96 #define yytos(stk) yystk_item( stk, 0 ) /* Evaluates to top-of-stack item. */
97
98 yystk_dcl( Yy_stack, int, YYMAXDEPTH );
99
100 typedef struct /* Typedef for value-stack elements. */
101 {
102     YYSTYPE left; /* Holds value of left-hand side attribute. */
103     YYSTYPE right; /* Holds value of current-symbol's attribute. */
104
105 } yyvstype;
106
107 yyvstype Yy_vstack[ YYMAXDEPTH ]; /* Value stack. */
108 yyvstype *Yy_vsp = Yy_vstack + YYMAXDEPTH; /* Value-stack pointer. */
109

```



Listing 4.4. continued...

```

110  @
111  @ Tables go here. LLama removes all lines that begin with @ when it copies
112  @                               llama.par to the output file.
113  ^L

```

4.5 are identical in content to the ones pictured in Figure 4.6 on page 210. Note that the `Yyd` table on lines 179 to 184 is not compressed because this output file was generated with the `-f` switch active. Were `-f` not specified, the tables would be pair compressed, as is described in Chapter Two. The `yy_act()` subroutine on lines 199 to 234 contains the `switch` that holds the action code. Note that \$ references have been translated to explicit value-stack references (`Yy_vsp` is the value-stack pointer). The `Yy_synch` array on lines 243 to 248 is a -1-terminated array of the synchronization tokens specified in the `%synch` directive.

Listing 4.5. llout.c—Parse Tables

```

114  /-----
115  * The YypNN arrays hold the right-hand sides of the various productions, listed
116  * back to front (so that they will be pushed in reverse order), NN is the
117  * production number (to be found in the symbol-table listing output with a -s
118  * command-line switch).
119  *
120  * Yy_pushtab[] is indexed by production number and points to the appropriate
121  * right-hand-side (YypNN) array.
122  */
123
124 YYPRIVATE int Yyp07[]={ 0 };
125 YYPRIVATE int Yyp06[]={ 261, 517, 260, 516, 2, 0 };
126 YYPRIVATE int Yyp09[]={ 5, 257, 4, 0 };
127 YYPRIVATE int Yyp08[]={ 518, 3, 0 };
128 YYPRIVATE int Yyp04[]={ 0 };
129 YYPRIVATE int Yyp03[]={ 259, 515, 258, 514, 1, 0 };
130 YYPRIVATE int Yyp05[]={ 261, 260, 0 };
131 YYPRIVATE int Yyp02[]={ 259, 258, 0 };
132 YYPRIVATE int Yyp01[]={ 256, 6, 513, 257, 512, 0 };
133 YYPRIVATE int Yyp00[]={ 0 };
134
135 YYPRIVATE int *Yy_pushtab[] =
136 {
137     Yyp00,
138     Yyp01,
139     Yyp02,
140     Yyp03,
141     Yyp04,
142     Yyp05,
143     Yyp06,
144     Yyp07,
145     Yyp08,
146     Yyp09
147 };
148
149 /-----
150 * Yyd[][] is the DFA transition table for the parser. It is indexed
151 * as follows:

```



Listing 4.5. continued...

```

152      *           Input symbol
153      *           +-----+
154      *           L | Production number   |
155      *           H |       or YYF        |
156      *           S |                   |
157      *           +-----+
158      *
159      * The production number is used as an index into Yy_pushtab, which
160      * looks like this:
161      *
162      *           Yy_pushtab      YypDD:
163      *           +-----+      +-----+
164      *           | *---|---->|           |
165      *           +-----+      +-----+
166      *           | *---|---->
167      *           +-----+
168      *           | *---|---->
169      *           +-----+
170      *
171      * YypDD is the tokenized right-hand side of the production.
172      * Generate a symbol table listing with llama's -l command-line
173      * switch to get both production numbers and the meanings of the
174      * YypDD string contents.
175      */
176
177 YYPRIVATE YY_TTYPE Yyd[ 6 ][ 7 ] =
178 {
179 /* 00 */ { 0, -1, -1, 1, 1, -1, -1 },
180 /* 01 */ { -1, -1, -1, 2, 2, -1, -1 },
181 /* 02 */ { -1, -1, -1, 5, 5, -1, -1 },
182 /* 03 */ { -1, 3, -1, -1, -1, 4, 4 },
183 /* 04 */ { -1, -1, -1, 8, 9, -1, -1 },
184 /* 05 */ { -1, 7, 6, -1, -1, 7, 7 }
185 };
186
187 /-----
188 * yy_next(state,c) is given the current state and input
189 * character and evaluates to the next state.
190 */
191
192 #define yy_next(state, c) Yyd[ state ][ c ]
193
194 /-----
195 * Yy_act() is the action subroutine. It is passed the tokenized value
196 * of an action and executes the corresponding code.
197 */
198
199 YYPRIVATE int yy_act( actnum )
200 {
201     /* The actions. Returns 0 normally but a nonzero error code can be returned
202     * if one of the acts causes the parser to terminate abnormally.
203     */
204
205     switch( actnum )
206     {
207         case 512:
208             {(Yy_vsp[1].right)=(Yy_vsp[2].right)=newname(); }
209             break;

```



Listing 4.5. continued...

```

210     case 513:
211         { freename((Yy_vsp[0].right)); }
212         break;
213     case 514:
214         { (Yy_vsp[1].right)=(Yy_vsp[2].right)=newname(); }
215         break;
216     case 515:
217         { yycode ("%s+=%s\n", Yy_vsp->left, (Yy_vsp[0].right));
218             freename((Yy_vsp[0].right)); }
219         break;
220     case 516:
221         { (Yy_vsp[1].right)=(Yy_vsp[2].right)=newname(); }
222         break;
223     case 517:
224         { yycode ("%s*=%s\n", Yy_vsp->left, (Yy_vsp[0].right));
225             freename((Yy_vsp[0].right)); }
226         break;
227     case 518:
228         { yycode ("%s=%0.*s\n", Yy_vsp->left, yylen, yytext); }
229         break;
230     default: printf("INTERNAL ERROR: Illegal act number (%s)\n", actnum);
231         break;
232     }
233     return 0;
234 }
235
236 /*-----
237 * Yy_synch[] is an array of synchronization tokens. When an error is
238 * detected, stack items are popped until one of the tokens in this
239 * array is encountered. The input is then read until the same item is
240 * found. Then parsing continues.
241 */
242
243 YYPRIVATE int Yy_synch[] =
244 {
245     RP,
246     SEMI,
247     -1
248 };
249
250 /*-----
251 * Yy_stok[] is used for debugging and error messages. It is indexed
252 * by the internal value used for a token (as used for a column index in
253 * the transition matrix) and evaluates to a string naming that token.
254 */
255
256 char *Yy_stok[] =
257 {
258     /* 0 */    "_EOI_",
259     /* 1 */    "PLUS",
260     /* 2 */    "TIMES",
261     /* 3 */    "NUM_OR_ID",
262     /* 4 */    "LP",
263     /* 5 */    "RP",
264     /* 6 */    "SEMI"
265 };
266

```



Listing 4.5. continued...

```

267
268 #ifdef YYDEBUG
269
270 /*-----
271 * Yy_snonterm[] is used only for debugging. It is indexed by the
272 * tokenized left-hand side (as used for a row index in Yyd[]) and
273 * evaluates to a string naming that left-hand side.
274 */
275
276 char *Yy_snonterm[] =
277 {
278     /* 256 */ "stmt",
279     /* 257 */ "expr",
280     /* 258 */ "term",
281     /* 259 */ "expr''",
282     /* 260 */ "factor",
283     /* 261 */ "term'"
284 };
285
286 /*-----
287 * Yy_sact[] is also used only for debugging. It is indexed by the
288 * internal value used for an action symbol and evaluates to a string
289 * naming that token symbol.
290 */
291
292 char *Yy_sact[] =
293 {
294     "{0}", "{1}", "{2}", "{3}", "{4}", "{5}", "{6}"
295 };
296
297 #endif

```

The arrays following line 256 provide human-readable diagnostics; they translate the various numeric values of the symbols into strings representing those symbols. *Yy_stok*, on lines 256 to 265, is indexed by token value and evaluates to a string representing the token's name. Similarly, *Yy_snonterm*[] on line 276 translates non-terminals, and *Yy_sact*[] on line 431 puts the actions in some sort of readable form. The {0}, and so forth, are mapped to the actual code in the *llout.sym* file, generated by giving **LLama** a *-s*, *-S*, or *-D* command-line switch.

Listing 4.6 starts moving into the parser proper. The streams declared on lines 300 to 302 are used for output to the code, bss, and data segments respectively. They're all initialized to *stdout*, but you can change them with an *fopen()* call.

The remainder of Listing 4.6 provides alternate versions of functions and macros for debugging (*YYDEBUG* defined) production modes. The macros and subroutines on lines 306 to 431 are used if debug mode is active, otherwise the same macros and subroutines are redefined on lines 431 to 497 not to print diagnostics, and so forth.

If *YYDEBUG* is defined, a second parse stack (*Yy_dstack*) is defined on line 308. This second, *symbol* stack exactly parallels the normal parse stack, but it holds strings representing the symbols, which are in turn represented by the numbers on the normal parse stack. For example, an *expr* is represented by the number 257 on the parse stack, and every time a 257 is pushed on the normal parse stack, the string "expr" is pushed onto the debugging stack. A string is popped (and discarded) from the symbol stack when a number is popped from the normal parse stack. This symbol stack is used in the debugging environment to make stack activity a little easier to follow.

Symbol-to-string conversion arrays: *Yy_stok*, *Yy_snonterm*, *Yy_sact*.

Output streams.

Debugging functions and macros, *YYDEBUG*.

Symbol stack, *Yy_dstack*.

Listing 4.6. *llama.par*— Macros for Parsing and Debugging Support

```

298 /*-----*/
299
300 FILE *yycodeout = stdout; /* Output stream for code. */
301 FILE *yybsfout = stdout; /* Output stream for initialized data. */
302 FILE *yydataout = stdout; /* Output stream for uninitialized data. */
303 int yynerrs = 0; /* Error count. */
304
305 /*-----*/
306 #ifdef YYDEBUG /* Debugging parse stack. */
307
308 yystk_dcl( Yy_dstack, char *, YYMAXDEPTH );
309
310 YYPRIVATE char *yy_sym( sym )
311 {
312     /* Return a pointer to a string representing the symbol, using the
313      * appropriate table to map the symbol to a string.
314      */
315
316     return ( YY_ISTERM( sym ) || !sym ) ? Yy_stok [sym]:
317         ( YY_ISNONTERM( sym ) ) ? Yy_snonterm [sym - YY_MINNONTERM]:
318             /* assume it's an act */ Yy_sact [sym - YY_MINACT];
319 }
320
321 /* Stack-maintenance. yy_push and yy_pop push and pop items from both the
322  * parse and symbol stacks simultaneously. The yycomment calls in both routines
323  * print a message to the comment window saying what just happened. The
324  * yy_pstack() call refreshes the stack window and also requests a new command
325  * from the user. That is, in most situations, yy_pstack() won't return until
326  * the user has typed another command (exceptions are go mode, and so forth).
327  * yy_pstack() also checks for break points and halts the parse if a breakpoint
328  * condition is met.
329 */
330
331 YYPRIVATE yy_push(x, val)
332 int x; /* Push this onto the state stack.*/
333 YYSTYPE val; /* Push this onto the value stack.*/
334 {
335     yypush ( Yy_stack, x );
336     yypush_( Yy_dstack, yy_sym(x) );
337
338     --Yy_vsp; /* The push() macro checked */
339     Yy_vsp->left = Yy_vsp->right = val; /* for overflow already. */
340
341     yycomment( "push %s\n", yy_sym( x ) );
342     yy_pstack( 0, 1 );
343 }
344
345 YYPRIVATE yy_pop()
346 {
347     int prev_tos = yypop( Yy_stack );
348     ++Yy_vsp;
349
350     yycomment( "pop %s\n", yypop_( Yy_dstack ) );
351     yy_pstack ( 0, 1 );
352
353     return prev_tos;
354 }
355

```



Listing 4.6. continued...

```

356 YYPRIVATE yy_say_whats_happening(tos_item, production)
357
358 int tos_item;           /* Item at top of stack */          */
359 int production;        /* production number we're about to apply */
360 {
361     /* Print a message in the comment window describing the replace operation
362      * about to be performed. The main problem is that you must assemble a
363      * string that represents the right-hand side of the indicated production.
364      * I do this using the appropriate Yy_pushtab element, but go through the
365      * array backwards (the Yy_pushtab arrays have been reversed to make the
366      * production-mode parse more efficient--you need to unreverse them here).
367      */
368
369     char buf[80];           /* Assemble string representing right-hand side here.*/
370     int count;             /* Maximum size of string representing RHS. */
371     int *start, *end;       /* Start and end of Yy_pushtab array that holds RHS. */
372
373     for( start = end = Yy_pushtab[ production ]; *end; end++ )    /* Find end.*/
374     ;
375
376     *buf = '\0';
377     for( count = sizeof(buf); --end >= start && count > 0 ; )          /* Assemble */
378     {                                                               /* string. */
379         strncat( buf, yy_sym(*end), count );
380         if( (count -= strlen( yy_sym(*end) + 1 )) < 1 )
381             break;
382         strncat( buf, " ", --count );
383     }
384
385     yycomment( "Applying %s->%s\n", yy_sym(tos_item), buf );
386 }
387
388 /* Use the following routines just like printf() to create output. In debug
389  * mode, all three routines print to the output window (yy_output() is in
390  * yydebug.c). In production mode, equivalent routines print to the associated
391  * streams (yycodeout, yybssout, or yydataout). The first argument to
392  * yy_output() tells the routine which stream is being used. If the stream is
393  * still set to the default stdout, then the message is written only to the
394  * window. If the stream has been changed, however, the output is sent both
395  * to the window and the associated stream.
396 */
397
398 yycode( fmt )                  /* Write something to the code-segment stream. */
399 char *fmt;
400 {
401     va_list args;
402     va_start( args, fmt );
403     yy_output( 0, fmt, args );
404 }
405
406 yydata( fmt )                 /* Write something to the data-segment stream. */
407 char *fmt;
408 {
409     va_list args;
410     va_start( args, fmt );
411     yy_output( 1, fmt, args );
412 }
413

```



Listing 4.6. continued...

```

414 yybss( fmt )                      /* Write something to the bss-segment stream. */
415     char *fmt;
416 {
417     va_list args;
418     va_start( args, fmt );
419     yy_output( 2, fmt, args );
420 }
421
422 /* Debugging versions of yycomment() and yy_error() are pulled out of yydebug.c
423 * when YYDEBUG is defined. Similarly, yy_break(), which halts the parse if a
424 * break-on-production-applied breakpoint has been triggered, is found in
425 * yydebug.c. It is eliminated from the production-mode output by defining it as
426 * an empty macro, below. Finally, yy_nextoken(), which evaluates to a yylex()
427 * call in production mode, and which is defined in yydebug.c, both gets the
428 * next token and prints it in the TOKEN window.
429 */
430
431 #else /*----- */
432
433 # define yy_push(x,v)      ( yypush(Yy_stack, x), \
434                                --Yy_vsp, Yy_vsp->left=Yy_vsp->right=v )
435
436 # define yy_pop()          ( ++Yy_vsp, yypop(Yy_stack) )
437 # define yy_nextoken()     yylex()
438 # define yy_quit_debug()   yy_sym()
439 # define yy_say_whats_happening(tos_item,prod)
440 # define yy_redraw_stack()
441 # define yy_pstack(refresh,print_it)
442 # define yy_break(x)
443
444
445 #ifndef va_list
446 #  include <stdarg.h>
447 #else
448 #  ifdef va_dcl
449     MUST USE ANSI VARIABLE-ARGUMENT CONVENTIONS IN <stdarg.h>
450 #  endif
451 #endif
452
453 void yycode( fmt, ... )
454 char *fmt;
455 {
456     va_list args;
457     va_start( args, fmt );
458     vfprintf( yycodeout, fmt, args );
459 }
460
461 void yydata( fmt, ... )
462 char *fmt;
463 {
464     va_list args;
465     va_start( args, fmt );
466     vfprintf( yydataout, fmt, args );
467 }
468
469 void yybss( fmt, ... )
470 char *fmt;
471 {
472     va_list args;
473     va_start( args, fmt );

```

Listing 4.6. continued..

```

474         vfprintf( yybssout, fmt, args );
475     }
476
477     void    yycomment( fmt, ... )
478     char    *fmt;
479     {
480         va_list args;
481         va_start( args, fmt );
482         vfprintf( stdout, fmt, args );
483     }
484
485     void    yyerror( fmt, ... )
486     char    *fmt;
487     {
488         va_list args;
489         extern char *yytext;
490         extern int yylineno;
491
492         va_start( args, fmt );
493         fprintf ( stderr, "ERROR (line %d near %s): ", yylineno, yytext );
494         vfprintf( stderr, fmt, args );
495         fprintf ( stderr, "\n" );
496     }
497 #endif

```

The symbol stack is manipulated in the push() and pop() subroutines on lines 331 to 354. Stack boundary checking is done only on the actual parse stack because, since all stacks are the same size and they all run in parallel, additional checking is redundant. These subroutines are translated into macros on lines 433 to 436 if YYDEBUG isn't active. Most other details are called out in comments in the listing, to which you are referred. I'll look at *yydebug.c*, which holds all the debug-mode support routines, in a moment.

The rest of *llout.c* is the error-recovery code and the actual parser, in Listing 4.7. Error-recovery is done with two subroutines on lines 502 to 545. *yy_in_synch()* (on line 502) checks to see if the symbol passed to it is in the set of synchronization symbols given to the %synch directive and listed in the *Yy_synch* table. *yy_synch()* does the actual error recovery. The error-recovery algorithm is described in the comment on line 516.

Listing 4.7. *llama.par*— The Parser

```

498  /*****
499   *  ERROR RECOVERY:
500   */
501
502 YYPRIVATE yy_in_synch( sym )
503 {
504     /* Return 1 if sym is in the synchronization set defined in Yy_synch. */
505
506     int *p ;
507
508     for( p = Yy_synch; *p &&  *p > 0 ; p++ )
509         if( *p == sym )
510             return 1;
511     return 0;
512 }

```



Listing 4.7. continued...

```

513
514 YYPRIVATE yy_synch( lookahead )
515 {
516     /* Recover from an error by trying to synchronize the input stream and the
517      * stack. Return the next lookahead token or 0 if we can't recover. Yyparse()
518      * terminates if none of the synchronization symbols are on the stack. The
519      * following algorithm is used:
520      *
521      * (1) Pop symbols off the stack until you find one in the synchronization
522      *      set.
523      * (2) If no such symbol is found, you can't recover from the error. Return
524      *      an error condition.
525      * (3) Otherwise, advance the input symbol either until you find one that
526      *      matches the stack symbol uncovered in (1) or you reach end of file.
527      */
528
529     int tok;
530
531     if( ++yynerrs > YYMAXERR )
532         return 0;
533
534     while( !yy_in_synch( tok = yytos( Yy_stack ) ) \
535             && !yystk_empty( Yy_stack ) ) /* 1 */
536         yy_pop();
537
538     if( yystk_empty( Yy_stack ) ) /* 2 */
539         return 0;
540
541     while( lookahead && lookahead != tok ) /* 3 */
542         lookahead = yy_nextoken();
543
544     return lookahead;
545 }
546
547 /-----
548 * The actual parser. Returns 0 normally, -1 if it can't synchronize after an
549 * error, otherwise returns a nonzero value returned by one of the actions.
550 */
551
552 int yyparse()
553 {
554     int *p; /* General-purpose pointer. */
555     YY_TTYPE prod; /* Production being processed. */
556     int lookahead; /* Lookahead token. */
557     int errcode = 0; /* Error code returned by an act. */
558     int tchar; /* Holds terminal character in yytext. */
559     int actual_lineno; /* Actual input line number. */
560     char *actual_text; /* Text of current lexeme. */
561     int actual_leng; /* Length of current lexeme. */
562     YYSTYPE val; /* Holds $$ value for replaced nonterm. */
563
564 # ifdef YYDEBUG
565     if( !yy_init_debug( Yy_stack, &yystk_p(Yy_stack),
566                         Yy_dstack, &yystk_p(Yy_dstack),
567                         Yy_vstack, sizeof(yyvstype), YYMAXDEPTH ) )
568         YYABORT;
569

```



Listing 4.7. continued...

```

570         yystk_clear(Yy_dstack);
571 # endif
572
573     yystk_clear(Yy_stack);
574     Yy_vsp = Yy_vstack + YYMAXDEPTH;
575
576     yy_push( YY_START_STATE, (Yy_vsp-1)->left ); /* Push start state onto */
577                                         /* parse stack and junk */
578                                         /* onto the value stack. */
579                                         /* User-supplied init. */
580     yy_init_llama( Yy_vsp );
581     lookahead = yy_nextoken();
582
583     while( !yystk_empty(Yy_stack) )
584     {
585         if( YY_ISACT( yytos(Yy_stack) ) )           /* if TOS is an action, do */
586         {                                         /* it and pop the action. */
587             yylineno      = ii_plineno() ;
588             yytext        = ii_ptext() ;
589             tchar         = yytext[ yyleng = ii_plength() ];
590             yytext[yyleng] = '\0' ;
591
592             if( errcode = yy_act( yytos(Yy_stack) ) )
593                 return errcode;
594
595             yy_pop();
596             yy_redraw_stack();
597             yytext[ yyleng ] = tchar ;
598         }
599         else if( YY_ISTERM( yytos(Yy_stack) ) ) /* Advance if it's a terminal.*/
600         {
601             if( yytos(Yy_stack) != lookahead ) /* ERROR if it's not there. */
602             {
603                 yyerror( "%s expected\n", Yy_stok[ yytos(Yy_stack) ] );
604                 if( !(lookahead = yy_synch(lookahead)) )
605                     YYABORT;
606             }
607             else
608             {
609                 /* Pop the terminal symbol at top of stack. Mark the current
610                  * token as the previous one (we'll use the previous one as
611                  * yytext in subsequent actions), and advance.
612                  */
613
614                 yy_pop();
615                 ii_mark_prev();
616
617                 lookahead      = yy_nextoken();
618                 actual_lineno = yylineno;
619                 actual_text   = yytext ;
620                 actual_leng   = yyleng ;
621             }
622             else
623             {
624                 /* Replace a nonterminal at top of stack with its right-hand side.
625                  * First look up the production number in the table with the
626                  * yy_next call. If prod==YYF, there was no legal transition and
627                  * error-processing is activated. Otherwise the replace operation
628                  * is done by popping the nonterminal, and pushing the right-hand
629                  * side from the appropriate Yy_pushtab entry. */
630             }
631         }
632     }
633 
```



Listing 4.7. continued...

```

630             */
631
632         prod = yy_next( yytos(Yy_stack)-YY_MINNTERM, lookahead );
633
634         if( prod == YYF )
635         {
636             yyerror( "Unexpected %s\n", Yy_stok[ lookahead ] );
637             if( !(lookahead = yy_synch(lookahead)) )
638                 YYABORT;
639         }
640         else
641         {
642             yy_say_whats_happening( yytos(Yy_stack), prod );
643             yy_break( prod );
644
645             val = Yy_vsp->right ;
646             yy_pop();
647
648             for( p = Yy_pushtab[ prod ] ; *p ; ++p )
649                 yy_push( *p, val );
650
651         }
652     }
653
654     yylineno = actual_lineno ; /* Make these hold reasonable values in case */
655     yytext   = actual_text   ; /* the remainder of the input file must be */
656     yyleng   = actual_leng   ; /* processed further. */                      */
657
658     yy_quit_debug();
659     YYACCEPT;
660 }
```

Top-down parser:
yyparse().

Phase problems with
yytext.

The actual parser, `yyvsparse()`, starts on line 552 of Listing 4.7. For the most part, it is a straightforward implementation of the algorithm in Table 4.12 on page 211. The one difficulty not covered in the algorithm is the lexeme, `yytext`, which must be valid when the action is executed. The problem is that the match-and-advance operation that's triggered when a lexeme is on top of the stack overwrites `yytext` with the new lexeme. Consequently, you have to mark the current lexeme as the previous one before advancing, using the `ii_mark_prev()` call on line 614 (which was put into the input routines described in Chapter Three for this very purpose). When a subsequent action is actually performed (starting on line 586), `yytext` is modified to reference the previous lexeme, not the current one—it references the lexeme associated with the token that we just advanced past, not the lookahead token. The code on lines 588, 589, and 596 is just null terminating the lexeme.

The remainder of `llout.c` is just copied from the input file at the end of Appendix D. I won't reproduce it here.

4.9.2 Occs and LLama Debugging Support—`yydebug.c`

This section discusses the debug-mode support routines used by the **LLama**-generated parser in the previous section. The same routines are used by the **occ**s-generated parser discussed in the next chapter. You should be familiar with the interface to the *curses*, window-management functions described in Appendix A before continuing.

The debugging module, yydebug.c, starts in Listing 4.8 with various definitions, which you should read now (there are extensive comments that describe what the variables and definitions do—I'll discuss some of them a bit more in a moment).

Listing 4.8. yydebug.c— Definitions

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <sys/types.h> /* ANSI/UNIX time functions. */
4 #include <sys/timeb.h> /* ANSI/UNIX time functions. */
5 #include <curses.h> /* Window functions. */
6 #include <signal.h> /* Needed by signal. */
7 #include <stdarg.h> /* ANSI variable-argument lists. */
8 #include <io.h> /* Prototypes for access(). */
9 #include <string.h> /* Prototypes for string functions. */
10 #include <stdlib.h> /* Prototypes for other library functions. */
11 #include <tools/debug.h> /* Various macros. */
12 #include <tools/l.h> /* Prototypes for all of l.lib, including all */
13 /* functions in the current file. */
14 extern char *yytext; /* Generated by LeX and lex. */
15 extern int yylineno;
16 extern int yyleng;
17
18 /* If your system doesn't have an <stdarg.h>, use the following:
19 *
20 * typedef char *va_list;
21 * #define va_start(arg_ptr,first) arg_ptr = (va_list)&first + sizeof(first)
22 * #define va_arg(arg_ptr,type) ((type*)(arg_ptr += sizeof(type)))[-1]
23 * #define va_end()
24 */
25 /* -----
26 * The following macros take care of system dependencies. They assume a 25-line
27 * screen on the IBM and a 24-line screen under Unix. Code put inside an MS()
28 * macro compiles only if MSDOS is #defined. Code in a UX() macro compiles only
29 * if MSDOS is not #defined. The NEWLINE define takes care of a bug in the UNIX
30 * curses package that isn't present in the DOS version presented in this book
31 * (it clears the bottom line after a scroll). box.h (in Appendix A) holds
32 * #defines for the IBM Box-drawing characters. #define NOT_IBM_PC to use the
33 * more portable defines in that file ('+' is used for corners, '|' for vertical
34 * lines, and '-' for horizontal ones) rather than the less-portable IBM
35 * graphics characters. fcntl() is also used only in UNIX mode.
36 */
37 #ifdef MSDOS
38 # include <tools/box.h>
39 # define SCRNSIZE 25
40 # define NEWLINE(win) (Interactive ? waddch( win, '\n') :0)
41 #else
42 # define NOT_IBM_PC
43 # include <tools/box.h>
44 # include <fcntl.h>
45 # define SCRNSIZE 24
46 # define NEWLINE(win) (Interactive ? (waddch( win, '\n'), wclrtoeol(win))\
47 : 0 )
48#endif
49
50 /* -----
51 * Defines for the windows. STACK_TOP is the top line of the stack window.
52 * DEFSTACK is the default size of the text area of the stack window.
53 * STACK_WINSIZE is the height of the stack window, including the border. IO_TOP
54 * is the top line of both the I/O and comments windows, and IO_WINSIZE is the
55 * height of the text area of those windows. It should use the whole screen

```



Listing 4.8. continued...

```

56     * less the area used for the stack and prompt windows.
57     */
58
59 #define STACK_TOP      0
60 #define DEFSTACK      11          /* Stacksize=DEFSTACK by default. */
61 #define STACK_WINSIZE (Stacksize +2)
62 #define PROMPT_TOP    (SCRNSIZE - 3)
63 #define PROMPT_WINSIZE 3
64 #define IO_TOP        (STACK_WINSIZE-1)
65 #define IO_WINSIZE   ((SCRNSIZE - (STACK_WINSIZE + PROMPT_WINSIZE)) + 2)
66
67 #define TOKEN_WIDTH   22         /* Width of token window including border. */
68 #define PRINTWIDTH    79         /* Only this many characters are printed on each
69                                * line by the write-screen (w) command. Extra
70                                * characters are truncated.
71                                */
72
73 #define ESC      0x1b          /* ASCII ESC character. */ *
74
75 /* -----
76  * Breakpoints. A breakpoint is set with a 'b' command. It causes automatic-mode
77  * operation to terminate immediately before applying a specific production or
78  * when a specified symbol is on the top of stack. P_breakpoint holds the
79  * production breakpoint; T_breakpoint holds the top-of-stack breakpoint;
80  * I_breakpoint is the input breakpoint. The former is an int because it's
81  * always a number. The latter two are strings because they can be symbolic
82  * names as well as numbers. The last variable, L_breakpoint, is the input-
83  * line breakpoint.
84 */
85
86 #define BRKLEN 33           /* Longest lexeme recognized in a breakpoint + 1. */
87
88 PRIVATE int    P_breakpoint = -1 ;
89 PRIVATE int    L_breakpoint = -1 ;
90 PRIVATE char   S_breakpoint[ BRKLEN ] = { '\0' } ;
91 PRIVATE char   I_breakpoint[ BRKLEN ] = { '\0' } ;
92
93 /*-----
94  * I've attempted to isolate these routines as much as possible from the actual
95  * parser. They do need to know where all the stacks are, however. The following
96  * variables are initialized at run-time by an access routine [yy_init_debug()]
97  * and are used to access static variables in the parser itself. Note that the
98  * addresses of the stack pointers are stored, not the contents of the stack
99  * pointers.
100 */
101
102 PRIVATE int    Abort      ;          /* Force input routine to return EOI. */
103 PRIVATE char   *Vstack    ;          /* Base address of value stack (or NULL */
104                                         /* if called by llama-generated parser).*/
105 PRIVATE int    Vsize      ;          /* Size of one element of value stack. */
106 PRIVATE char   **Dstack   ;          /* Base address of debug (symbol) stack.*/
107 PRIVATE char   ***P_dsp   ;          /* Pointer to debug-stack pointer. */
108 PRIVATE int    *Sstack    ;          /* Base address of state stack. */
109 PRIVATE int    **P_sp    ;          /* Pointer to state-stack pointer. */
110 PRIVATE int    Depth      ;          /* Stack depth (all three stacks). */
111
112 /*-----
113  * The following variables are all used internally
114 */
115

```

Listing 4.8. continued...

```

116 PRIVATE WINDOW *Stack_window ; /* Windows for the debugging screen. */
117 PRIVATE WINDOW *Prompt_window ;
118 PRIVATE WINDOW *Code_window ;
119 PRIVATE WINDOW *Comment_window ;
120 PRIVATE WINDOW *Token_window ;
121 PRIVATE int Stacksize = DEFSTACK; /* Number of active lines in the stack */
122 /* window (doesn't include border). */
123 PRIVATE int Onumele = 0; /* Number of elements on the stack. */
124 PRIVATE int Interactive = 1; /* Interactive mode (not n or N). */
125 PRIVATE int Singlestep = 1; /* Single step through parse if true. */
126 PRIVATE long Delay = 0L; /* Amount of time to wait after printing */
127 /* each stack update when not single */
128 /* stepping (milliseconds). */
129 PRIVATE int Inp_fm_file = 0; /* 1 if input file is open. */
130 PRIVATE FILE *Log = NULL; /* Pointer to the log file if one is open. */
131 PRIVATE int No_comment_pix = 0; /* 1 if no comment-window output is printed. */
132 PRIVATE int No_stack_pix = 0; /* 1 if no stack pictures are to be printed */
133 /* in the log file. */
134 PRIVATE int Horiz_stack_pix = 0; /* 1 if stack pictures are printed horiz- */
135 /* entally in the log file. */
136 PRIVATE int Parse_pix; /* if(Horiz_stack_pix), print state stack. */
137 PRIVATE int Sym_pix; /* if(Horiz_stack_pix), print symbol stack. */
138 PRIVATE int Attr_pix; /* if(Horiz_stack_pix), print attrib. stack. */
139
140 #ifndef MSDOS /*----- UNIX SYSTEM V ONLY -----*/
141 /* Since MS-DOS has a system call that
142 * gives the keyboard status, detecting
143 * if a key has been pressed without
144 * reading the character is easy. In
145 * UNIX you must use SIGIO to set a
146 * flag (Char_avail). kbready() is the
147 * SIGIO exception handler.
148 */
149 PRIVATE int Char_avail = 0;
150 #define kbhit() Char_avail
151
152 #else /*-----DOS VERSION ONLY-----*/
153
154 extern int kbhit ( void ); /* Microsoft function. returns 1 if a */
155 /* character is waiting to be read */
156 /* from the keyboard buffer. This */
157 /* function is provided in most */
158 /* MS-DOS compiler's libraries. */
159
160 /* The Map[] array converts IBM box-drawing characters to something that's
161 * printable. The corners are converted to plus signs, horizontal lines go to
162 * dashes, and vertical lines map to vertical bars. The conv() subroutine is
163 * passed a character and returns a mapped character. It must be a subroutine
164 * because of the way that it's used below. It would have unacceptable
165 * side-effects if rewritten as a macro.
166 */
167
168 PRIVATE unsigned char Map[] =
169 {
170     '|', '+', '+', '+', '+', '+', '+', '|', '+', '+', '+', '+', '+', '+',
171     '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+',
172     '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+', '+'
173 };
174

```



Listing 4.8. continued...

```

175 PRIVATE int conv(c)
176 {
177     return (VERT <= c && c <= UL) ? Map[c - VERT] : c ;
178 }
179 #endif /*-----*/

```

Initialize debugging:
yy_init_debug().

The actual code starts in Listing 4.9. The `yy_init_debug()` function on line 180 initializes the debugging functions. It is passed pointers to various variables in the parser proper that it needs to draw the windows (things like the base addresses of the stacks). Note that pointers to the stack pointers are passed, rather than the contents of the stack pointers. This way the parser can push and pop stuff at will without having to tell the debugging routines every time it does a push or pop. The debugging routines can just examine the pointers themselves. The `vstack` argument is `NULL` if a **LLama**-generated parser is active (because **LLama** doesn't use a value stack). In this case, the value stack is not printed in the log file or the stack window.

Getting keyboard status
under UNIX, SIGIO.

The `signal()` call on line 197 is used to determine keyboard status in UNIX systems. (The `UX` macro is defined in `debug.h`. It's discussed in Appendix A.) The `SIGIO` signal is generated every time a key is struck. Here, the `signal` activates the `kbready()` function (declared on line 257) which sets a global flag (`Char_avail`) to true. This flag can be examined when you're looking to see if a character is ready. The flag is explicitly set back to zero when the character is finally read. The `fcntl()` call on line 198 enables `SIGIO` (the signal is not generated otherwise).

Disable curses on abort.

The second `signal()` call on line 211 is used to disable curses when a Ctrl-Break (Ctrl-C or DEL in UNIX) comes along. Curses must be turned off explicitly in this case to prevent the screen from being left in an unknown state (with characters not echoing, and so forth). The handler is on line 263. Note that it disables `SIGINT` on line 265 so that another Ctrl-Break won't interrupt the shut-down process (this is not a problem on newer UNIX versions, but it can't hurt). The `q` command to the debugger quits by raising the `SIGINT` signal, which will execute the handler. This way, any cleaning up of temporary files that is done by a signal handler in the compiler itself is also done when you quit the debugger. You can use the `a` command if you want to leave the debugger without cleaning up.

`q` command raises
`SIGINT`.

Input buffering and character echo is turned off on line 213 of Listing 4.9. Note that this is necessary in the BSD curses that runs on the UNIX system that I use, but may be contraindicated on other systems. Ask your system administrator if this is the correct thing to do with your implementation. The windows themselves are opened on the next few lines. The `boxwin()` function (defined on line 287) works much like a standard `newwin()` call. It opens up two windows: the outer one as a subwindow to `stdscr`, and the inner one as a subwindow to the outer one. The outer window holds the box and the inner window holds the text. This way, text written to the inner window won't overwrite the box. The window title is printed, centered, in the top line of the box.

Parsing command-line ar-
guments:
`yy_get_args()`.

The final routine of interest in Listing 4.9 is the `yy_get_args()` function on line 320. This routine parses the command line for a stack size (specified with `-s`) and an input-file name. If you don't want to get this information from the command line, you can do the following:

```

floccinaucinihilipilification()
{
    static char *vects[] = /* Simulate argv. */ {
        "", /* Junk */
        "-s18", /* Stack-window size == 18 lines. */
        "test" /* Input file name. */
    };

    yy_get_args( 3, vects );
}

```

Listing 4.9. yydebug.c—Initialization

```

180 PUBLIC int yy_init_debug(sstack, p_sp, dstack, p_dsp, vstack, v_ele_size, depth)
181
182 int *sstack; /* Base address of state stack. */
183 int **p_sp; /* Address of state-stack pointer. */
184 char **dstack; /* Address of debug stack. */
185 char ***p_dsp; /* Address of debug-stack pointer. */
186 void *vstack; /* Base address of value stack or NULL for LLama. */
187 int v_ele_size; /* Size of one element of value stack. */
188 int depth; /* Number of elements in all three stacks. */
189 {
190     /* Initialize for interactive I/O for curses. Return 1 on a successful
191     * initialization, 0 otherwise.
192     */
193
194     char buf[80];
195
196     UX( int flags; )
197     UX( signal( SIGIO, kready ) );
198     UX( flags = fcntl( fileno(stdin), F_GETFL, 0 ) );
199     UX( fcntl( fileno(stdin), F_SETFL, flags | FASYNC ) );
200
201     Sstack = sstack;
202     Dstack = dstack;
203     Vstack = (char *) vstack;
204     Vsize = v_ele_size;
205     P_sp = p_sp;
206     P_dsp = p_dsp;
207     Depth = depth;
208     Abort = 0;
209
210     initscr();
211     signal( SIGINT, die_a_horrible_death );
212
213     noecho(); /* Don't echo input characters automatically. */
214     crmode(); /* Don't buffer input. */
215     MS( nosave()); /* Don't save region under windows (my curses only). */
216
217     Stack_window = boxwin( STACK_WINSIZE, 80, STACK_TOP, 0, "[stack]" );
218     Comment_window = boxwin( IO_WINSIZE, 40, IO_TOP, 0, "[comments]" );
219     Code_window = boxwin( IO_WINSIZE, 41, IO_TOP, 39, "[output]" );
220
221     Prompt_window = boxwin(PROMPT_WINSIZE, (80 - TOKEN_WIDTH) + 1,
222                             PROMPT_TOP, 0, "[prompts]" );
223
224     Token_window = boxwin(PROMPT_WINSIZE, TOKEN_WIDTH, PROMPT_TOP,
225                           80 - TOKEN_WIDTH, "[lookahead]" );
226     scrolllok( Stack_window, TRUE );

```

Listing 4.9. continued...

```

227     scrolllok( Comment_window, TRUE );
228     scrolllok( Code_window,    TRUE );
229     scrolllok( Prompt_window, TRUE );
230     scrolllok( Token_window,  TRUE );
231     wrapok( Token_window, FALSE );
232
233     Onumele = 0;
234
235     while( !Inp_fm_file )
236     {
237         /* If you don't have an input file yet, get one. yyprompt() prints the
238          * prompt in the PROMPT window, and fills buf with the reply.
239          */
240
241         if( !yyprompt( "Input file name or ESC to exit: ", buf, 1 ) )
242         {
243             yy_quit_debug();
244             return 0;
245         }
246         new_input_file( buf );
247     }
248     delay();           /* Wait for a command before proceeding.*/
249     return 1;
250 }
251
252 /-----
253 * Exception handlers:
254 */
255
256 #ifndef MSDOS
257 PRIVATE void      kbready()           /* Called when new character is available. */
258 {
259     Char_avail = 1;
260 }
261 #endif
262
263 PRIVATE void die_a_horrible_death()           /* Come here on a SIGINT */
264 {                                              /* or 'q' command. */
265     signal( SIGINT, SIG_IGN );
266     yy_quit_debug();
267     exit( 0 );
268 }
269
270 PUBLIC void yy_quit_debug()           /* Normal termination. */
271 {
272     echo();           /* Turn echo and editing back on. */
273     nocrmode();
274     move( 24, 0 );   /* Put the cursor on the bottom of the screen. */
275     refresh();
276     endwin();        /* Turn off curses. */
277
278     if( Log )
279         fclose( Log );
280
281     stop_prnt();
282     signal( SIGINT, SIG_DFL );
283 }
284

```



Listing 4.9. continued...

```

285  /*-----*/
286
287 PRIVATE WINDOW *boxwin( lines, cols, y_start, x_start, title )
288 int    lines;
289 int    cols;
290 int    y_start;
291 int    x_start;
292 char   *title;
293 {
294     /* This routine works just like the newwin() except that the window has a
295      * box around it that won't be destroyed by writes to the window. It
296      * accomplishes this feat by creating two windows, one inside the other,
297      * with a box drawn in the outer one. It prints the optional title centered
298      * on the top line of the box. Set title to NULL (or "") if you don't want
299      * a title. Note that all windows are made subwindows of the default window
300      * to facilitate the print-screen command.
301
302
303     WINDOW *outer;
304
305     outer = subwin( stdscr, lines, cols, y_start, x_start);
306     box( outer, VERT, HORIZ );
307
308     if( title && *title )
309     {
310         wmove( outer, 0, (cols - strlen(title))/2 );
311         wprintw( outer, "%s", title );
312     }
313
314     wrefresh( outer );
315     return subwin( outer, lines-2, cols-2, y_start+1, x_start+1 );
316 }
317
318 /*-----*/
319
320 PUBLIC int yy_get_args( argc, argv )
321 char **argv;
322 {
323     /* Scan argv arguments for the debugging options and remove the arguments
324      * from argv. Recognized arguments are:
325      *
326      * -sN      Set stack-window size to N lines. The size of the other windows
327      *           scale accordingly. The stack window is not permitted to get so
328      *           large that the other windows will disappear, however.
329      *
330      * The first argument that doesn't begin with a minus sign is taken to be
331      * the input file name. That name is not removed from argv. All other
332      * arguments are ignored and are not removed from argv, so you can process
333      * them in your own program. This routine prints an error message and
334      * terminates the program if it can't open the specified input file.
335      * Command-line processing stops immediately after the file name is
336      * processed. So, given the line:
337      *
338      *           program -x -s15 -y foo -s1 bar
339      *
340      * Argv is modified to:
341      *
342      *           program -x -y foo -s1 bar
343      *
344      * The file "foo" will have been opened for input and the stack window will

```



Listing 4.9. continued...

```

345      * be 15 lines high. Return new value of argc that reflects the removed
346      * arguments;
347      */
348
349      char    **newargv;
350      char    **oldargv = argv;
351      char    *filename = NULL;
352      int     ssize      = DEFSTACK;
353
354      newargv = ++argv;
355      for( --argc; --argc >= 0; ++argv )
356      {
357          if( argv[0][0] != '-' )
358          {
359              filename = *newargv++ = *argv;
360              break;
361          }
362          else if( argv[0][1] == 's' )           /* -s */
363              ssize = atoi( &argv[0][2] );        /* Don't copy to *newargv here. */
364          else                                /* -? */
365              *newargv++ = *argv;
366      }
367
368      Stacksize = ( ssize < 1                  ) ? DEFSTACK :           /* -? */
369      ( ssize > (SCRNSIZE-6) ) ? SCRNSIZE-6 :           /* -? */
370      /* ssize is in bounds */      ssize;
371
372      if( filename )
373      {
374          /* Open input file if one was specified on the command line. */
375
376          if( ii_newfile(filename) != -1 )
377              Inp_fm_file = 1;
378          else
379          {
380              perror( filename );
381              exit( 1 );
382          }
383      }
384      return newargv - oldargv;
385  }

```

Output functions: yyerror()
yyerror().

The next listing (Listing 4.10) holds all the output functions. There is one such routine for each window. In addition, yyerror() (on line 510), writes to the comment window and simulates the standard yyerror() function by adding an input line number and token value to the error message. In addition to the standard output functions, display_file() (on line 558 of Listing 4.10) is used by the f command to print a file in the stack window, and write_screen() is used by the w command to save the current screen to a file. Note that, since everything's a subwindow to stdscr, the stdscr functions can be used on line 654 to read the entire screen (rather than doing it window by window).

Listing 4.10. yydebug.c—Window Output Functions

```

386 PRIVATE void prnt_putc(c, win)
387 WINDOW *win;
388 {
389     /* All output done through prnt_putc is suppressed in Go mode. Also note
390      * that the arguments are reversed from addch(). This reversal lets you use
391      * the prnt() subroutine (described in Appendix A), which expects a putc()-like
392      * output function. Newlines are suppressed here so that you can have
393      * more control over scrolling. Similarly, sequences of white space are
394      * replaced by a single space character to conserve space in the window.
395      * Test_c is used to take care of the IBM graphics characters that form
396      * the vertical line separating the stream-identification column from the
397      * actual output. The c is mapped to a '|' if it's too large to be an ASCII
398      * character (so isspace() will work properly).
399 */
400
401 static WINDOW *last_win = NULL;
402 static int last_c = 0;
403 int test_c;
404
405 if( Interactive && c != '\n' )
406 {
407     test_c = (c < 0x7f) ? c : '|';
408
409     if( !(win==last_win && isspace(test_c) && isspace(last_c)) )
410         waddch( win, isspace(test_c) ? ' ' : c );
411
412     last_win = win;
413     last_c = test_c;
414 }
415 }
416
417 PRIVATE void refresh_win( win )
418 WINDOW *win;
419 {
420     /* Refresh the windows if you're not in go mode. (If you are, nothing will
421      * have been written, so there's not point in doing the refresh
422      */
423
424     if( Interactive )
425         wrefresh( win );
426 }
427
428 /*-----*/
429
430 PUBLIC void yy_output( where, fmt, args )           /* Generate code */
431 int where;
432 char *fmt;
433 va_list args;
434 {
435     /* Works like vprintf(), but sends output to the code window. In the window,
436      * it ignores any newlines in the string but prints a newline after every
437      * call. All code sent to yycode(), yydata(), and yybss() is funneled
438      * here. "where" should be one of the following:
439      *
440      *      0 : code
441      *      1 : data
442      *      2 : bss
443      *
444      * Note that if the three associated streams (yycodeout, yybssout, and

```



Listing 4.10. continued...

```

445     * yydataout--all declared in the parser output file) are not directed to
446     * stdout, output is sent to that stream TOO. Don't modify these to point
447     * at stderr (or any other stream that accesses the console: /dev/tty, con,
448     * etc.) or you'll mess up the screen.
449     *
450     * Note that the real yycode(), etc (ie. the ones supplied when YYDEBUG
451     * is not defined) don't do anything special with newlines. In particular,
452     * they are not inserted automatically at the end of the line. To make both
453     * sets of routines compatible, your output strings should all have exactly
454     * one newline, placed at the end of the string (don't imbed any in the
455     * middle).
456     */
457
458 extern FILE *yycodeout, *yydataout, *yybssout;
459
460 if( Log )
461 {
462     fprintf( Log,      where == 0 ? "CODE->" :
463             where == 1 ? "DATA->" : "BSS-->" );
464     prnt( fputc, Log, fmt, args );
465     fputc( '\n', Log );
466 }
467
468 NEWLINE( Code_window );
469
470 prnt_putc( where==0 ? 'C' : where==1 ? 'D' : 'B',      Code_window );
471 prnt_putc( VERT,                                     Code_window );
472
473 prnt( prnt_putc, Code_window, fmt, args );
474 refresh_win( Code_window );
475
476 if( where == 0 && yycodeout != stdout )
477     vfprintf( yycodeout, fmt, args );
478
479 if( where == 1 && yydataout != stdout )
480     vfprintf( yydataout, fmt, args );
481
482 if( where == 2 && yybssout != stdout )
483     vfprintf( yybssout, fmt, args );
484 }
485
486 /*-----*/
487
488 PUBLIC void yycomment( fmt, ... )
489 char      *fmt;
490 {
491     /* Works like printf() except that it automatically prints a newline
492     * IN FRONT OF the string and ignores any \n's in the string itself. Writes
493     * into the comment window, and outputs a message to the log file if
494     * logging is enabled.
495     */
496
497     va_list      args;
498     va_start( args, fmt );
499
500     if( Log && !No_comment_pix )
501         prnt( fputc, Log, fmt, args );
502

```



Listing 4.10. continued...

```

503     NEWLINE( Comment_window );
504     prnt( prnt_putc, Comment_window, fmt, args );
505     refresh_win( Comment_window );
506 }
507
508 /*-----*/
509
510 PUBLIC void yyerror( fmt, ... )
511 char *fmt;
512 {
513     /* Debugging version of the error routine. Works just like the nondebugging
514      * version, but writes to the comment window. Note that yycomment() copies
515      * the error message to the Log file. Interactive mode is temporarily
516      * enabled to assure that error messages get printed.
517 */
518
519     int old_interactive;
520     va_list args;
521     va_start( args, fmt );
522
523     old_interactive = Interactive;
524     Interactive = 1;
525
526     yycomment( "ERROR, line %d near <%s>\n", yylineno, yytext );
527
528     if( Log )
529         prnt( fputc, Log, fmt, args );
530
531     NEWLINE( Comment_window );
532     prnt( prnt_putc, Comment_window, fmt, args );
533     refresh_win( Comment_window );
534
535     Interactive = old_interactive;
536     Singlestep = 1; /* Force a breakpoint */
537     yy_pstack( 0, 1 );
538 }
539
540 /*-----*/
541
542 PRIVATE void yy_input( fmt, ... )
543 char *fmt;
544 {
545     /* This is not an input function; rather, it writes to the INPUT window.
546      * It works like printf(). Note that nothing is logged here. The logging
547      * is done in nextoken(). Ignores all \n's in the input string.
548 */
549
550     va_list args;
551     va_start( args, fmt );
552     prnt( prnt_putc, Prompt_window, fmt, args );
553     refresh_win( Prompt_window );
554 }
555
556 /*-----*/
557
558 PRIVATE void display_file( name, buf_size, print_lines )
559 char *name; /* Initially holds the file name, but */
560 int buf_size; /* recycled as an input buffer. */
561 int print_lines;
562 {
```



Listing 4.10. continued...

```

563     /* Display an arbitrary file in the stack window, one page at a time.
564      * The stack window is not refreshed by this routine.
565      */
566
567     FILE  *fd;
568     int    i;
569     int    lineno = 0;
570
571     if( !(fd = fopen( name, "r")) ) {
572     {
573         NEWLINE  ( Prompt_window );
574         wprintw ( Prompt_window, "Can't open %s", name );
575         wrefresh ( Prompt_window );
576         presskey ();
577     }
578     else
579     {
580         /* Note that order of evaluation is important in */
581         /* the following while statement. You don't want */
582         /* to get the line if i goes past 0. */
583
584         for( i = Stacksize-1 ;; i = (*name == ' ') ? 1 : Stacksize - 2 )
585         {
586             while( --i >= 0 && fgets(name, buf_size, fd) )
587             {
588                 if( print_lines )
589                     wprintw( Stack_window, "%3d:", ++lineno );
590
591                 wprintw ( Stack_window, "%s", name );
592                 wrefresh( Stack_window );
593             }
594
595             if( i > 0 )
596                 break;
597
598             if(!yyprompt("ESC quits. Space scrolls 1 line. Enter for screenful",
599                         name, 0) )
600                 break;
601             yyprompt("*** End of file. Press any key to continue ***", name, 0 );
602             fclose( fd );
603         }
604
605     /*-----*/
606
607     PRIVATE void write_screen( filename )
608     char   *filename;
609     {
610         /* Print the current screen contents to the indicated file. Note that the
611          * right edge of the box isn't printed in order to let us have 79-character
612          * lines. Otherwise, the saved screen shows up as double-spaced on most
613          * printers. The screen image is appended to the end of the file. In MS-DOS,
614          * Use "prn:" as the file name if you want to go to the printer.
615          *
616          * Syserrlist and errno are both defined in <stdlib.h>
617     */
618
619     char       buf[2];
620     char       *mode = "a";
621     int        row, col, y, x;
622     FILE      *file;

```

Listing 4.10. continued...

```

623
624     if( access( filename, 0 ) == 0 )
625         if( !yyprompt("File exists, overwrite or append? (o/a) : ", buf, 0 ) )
626     {
627         NEWLINE ( Prompt_window );
628         yy_input( "Aborting command." );
629         presskey();
630         return;
631     }
632     else
633     {
634         if( toupper(*buf) == 'O' )
635             mode = "w";
636     }
637
638     if( file = fopen( filename, mode ) )
639         yy_input("...%s %s...", 
640                  *mode=='w' ? "overwriting" : "appending to", filename );
641     else
642     {
643         yy_input( "Can't open %s: %s.", filename, sys_errlist[errno] );
644         presskey();
645         return;
646     }
647
648     getyx( Prompt_window, y, x );
649
650     for( row = 0; row < SCRNSIZE; row++ )
651     {
652         for( col = 0; col < PRINTWIDTH; col++ )
653         {
654             UX( fputc(      mvinch(row,col) , file ); )
655             MS( fputc( conv( mvinch(row,col) ), file ); )
656         }
657
658         fputc( '\n', file );
659     }
660
661     fclose( file );
662     wmove( Prompt_window, y, x );
663 }
```

The routines in Listing 4.11 do the real work. `yy_pstack()` (on line 664) is called by the parser every time the stack is modified. The stack is printed to the log file, if necessary, on line 700 to 773. `Interactive` is false on line 775 if the debugger is in noninteractive mode (initiated with an `n` command). In this case, a speedometer readout that tells you that the program is actually doing something is updated (it can take a while to parse a big input file, even in noninteractive mode) and the routine returns. A Stack breakpoint is triggered on line 785 if necessary (the debugger is just thrown back into single-step mode if one is found).

The `delay()` call on line 874 does one of two things. If you're not single stepping, it just delays by zero or more seconds (you can modify the number with the `d` command). `delay()` gets and executes a command if you're single stepping or you hit a key during the delay. `delay()` doesn't exit until one of the commands that starts up the parse again (space to singlestep, `n` to enter noninteractive mode, or `g` to go) is executed. The `delay()` subroutine itself starts at the top of the following Listing (Listing 4.12). Note

Update stack window,
`YY_pstack()`.

Stack breakpoint.

Main command loop,
`delay()`.

that I'm using the UNIX/ANSI time `ftime()` function here to get the time in milliseconds. It loads the `time_buf` structure with the elapsed number of seconds since January 1, 1970 (`time_buf.time`) and the number of milliseconds as well (in `time_buf.millitm`). The `Delay` variable holds the desired delay (in milliseconds).

Listing 4.11. *yydebug.c*—Stack-Window Maintenance and the Control Loop

```

664 PUBLIC void yy_pstack( do_refresh, print_it )
665
666 int do_refresh; /* redraw entire window rather than update */
667 int print_it; /* if true, print the stack to the log file */
668 {
669     /* Print the state, debug, and value stacks.
670      *
671      * The library routine yppstk() (which returns an empty string by default)
672      * is called to print value stack items. It should return a pointer to a
673      * string that represents the value-stack contents. Just provide a similarly
674      * named routine to print your real value stack. The LLAMA parser passes
675      * garbage to yppstk(), so use the default routine in LLAMA. The OCCS parser
676      * passes yppstk() a pointer to a value-stack item and it should return a
677      * pointer to a string representing attributes of interest. The line should
678      * not contain any newlines and it should be at most 58 characters long.
679      * If do_refresh is true, the entire stack window is redrawn, otherwise
680      * only those parts of the window that have been changed are modified.
681      */
682
683     int numele ; /* # of elements on the stack */ /
684     int *toss ; /* top of state stack */ /
685     char **tods ; /* top of debug stack */ /
686     char *tovs ; /* top of value stack */ /
687     int *state ; /* current state-stack pointer */ /
688     char **debug ; /* current debug-stack pointer */ /
689     char *value ; /* current value-stack pointer */ /
690     int width ; /* Width of column in horiz. stack */ /
691     static int times_called = -1 ; /* # of times this subroutine called */ /
692     char *p ;
693     int i ;
694
695     state = *P_sp;
696     debug = *P_dsp;
697     numele = Depth - (state - Sstack);
698     value = (Vstack + (Depth - numele) * Vsize);
699
700     if( Log && !No_stack_pix && print_it )
701     {
702         /* Print the stack contents out to the log file. */
703
704         if( !Horiz_stack_pix )
705         {
706             fprintf(Log, " +---+-----+\n");
707             if( numele <= 0 )
708                 fprintf( Log, " * | * | *****| Stack is empty.\n");
709             else
710             {
711                 toss = state;
712                 todः = debug;
713                 tovs = value;
714                 for(i = numele; --i >= 0; ++toss, ++todः, tovs += Vsize )
715                     fprintf( Log, "%3d|%3d| %16.16s | %1.52s\n",
716                             toss - state, *toss, *todः, yppstk(tovs, *todः));
717             }
718         }
719     }

```



Listing 4.11. continued...

```

718         fprintf(Log, " +---+-----+\n");
719     }
720     else
721     {
722         if( state < Sstack )
723             fprintf( Log,"*** Stack empty ***\n");
724         else
725         {
726             /* Print horizontal stack pictures. Note that you have to go
727              * through the stack from bottom to top to get the top-of-stack
728              * element on the right.
729             */
730
731             for( i = 0; i <= 2; ++i )
732             {
733                 if( !Parse_pix && i == 0 ) continue;
734                 if( !Sym_pix && i == 1 ) continue;
735                 if( !Attr_pix && i == 2 ) continue;
736
737                 switch( i )
738                 {
739                     case 0: fprintf( Log, " PARSE " ); break;
740                     case 1: fprintf( Log, " SYMBOL " ); break;
741                     case 2: fprintf( Log, " ATTRIB " ); break;
742                 }
743
744                 toss = Sstack + (Depth - 1);
745                 todः = Dstack + (Depth - 1);
746                 tovs = Vstack + ((Depth - 1) * Vsize);
747
748                 for(; toss >= state; --toss, --todः, tovs -= Vsize )
749                 {
750                     /* Find width of the column. I'm assuming that the
751                      * numbers on the state stack are at most 3 digits
752                      * long, if not, change the 3, below.
753                     */
754
755                     p      = yypstk( tovs, *todः );
756                     width = 3;
757
758                     if( Sym_pix ) width = max( width, strlen(*todः) );
759                     if( Attr_pix ) width = max( width, strlen(p) );
760
761                     switch( i )
762                     {
763                         case 0: fprintf( Log, "%-*d ", width, *toss ); break;
764                         case 1: fprintf( Log, "%-*s ", width, *todः ); break;
765                         case 2: fprintf( Log, "%-*s ", width, p ); break;
766                     }
767
768                     fputc( '\n', Log );
769                 }
770             }
771         }
772     }
773
774     if( !Interactive )
775     {

```



Listing 4.11. continued...

```

777 if( ++times_called % 25 == 0 )
778 {
779     wprintw ( Stack_window, "working: %d\r", times_called );
780     wrefresh( Stack_window );
781 }
782 return;
783 }
784
785 if( *S_breakpoint && state < Sstack + Depth )
786 {
787     /* Break if the breakpoint is a digit and the top-of-stack item has that
788      * value, or if the string matches the string currently at the top of
789      * the symbol stack.
790      */
791
792     if( isdigit(*S_breakpoint) )
793     {
794         if( atoi(S_breakpoint) == *state )
795             Singlestep = 1;
796     }
797     else if( !strcmp(S_breakpoint, *debug) )
798         Singlestep = 1;
799 }
800
801 if( do_refresh )
802     YY_redraw_stack();                                     /* Redraw entire stack */
803
804 else if( numele > Onumele )
805 {
806     /* The stack has grown. Redraw only those parts of the stack that have
807      * changed. (I'm assuming only by one element.) The main difficulty
808      * here is that only the top few elements of a large stack are
809      * displayed. Consequently, the stack window may have to scroll up
810      * or down a line if the stack size is hovering around the window size.
811      * There's no portable way to scroll the window up under UNIX curses, so
812      * we have to redraw the stack to scroll up in this situation. We'll
813      * overwrite the top element with itself by the wprintw() call, but
814      * that's no big deal, and it simplifies the code.
815      */
816
817     if( numele > Stacksize )           /* scroll down, opening up top line */
818     {
819         MS( wscroll( Stack_window, -1 );      )
820         UX( yy_redraw_stack();            )
821
822         wmove( Stack_window, 0, 0 );
823     }
824     else
825         wmove( Stack_window, Stacksize - numele, 0 );
826
827     wprintw( Stack_window, "%3d%c %16.16s %c %1.52s",
828             *state, VERT, *debug, VERT, yypstk(value, *debug) );
829     wrefresh( Stack_window );
830 }
831 else
832 {
833     /* The stack has shrunk, perhaps by several elements. Remove them one at
834      * a time. (It's too confusing if several elements disappear from the
835      * stack at once. It's best to watch them go one at a time.) If the
836      * number of elements on the stack (i) is greater than the window size,

```

Listing 4.11. continued...

```

837             * you can pop an element by scrolling up and then writing in a new
838             * bottom line, otherwise, just go to the correct line and erase it.
839             * Do a refresh after each pop.
840             */
841
842     for( i = Onumele; i > numele; --i )
843     {
844         if( i > Stacksize )
845         {
846             /* Do a pop by scrolling up, the easiest way to scroll is to
847             * move to the right edge of the bottom line and then issue
848             * a newline. After the scroll, overwrite the now-blank bottom
849             * line with the appropriate stack information. The involved
850             * expression that is the first argument to yypstk is doing:
851             *           (Vstack + Depth)[-i + Stacksize]
852             * It must do the pointer arithmetic explicitly, however, by
853             * multiplying by the size of one value-stack item (Vsize).
854             */
855
856             wmove( Stack_window, Stacksize-1, 77 );
857             NEWLINE( Stack_window );
858             wprintw( Stack_window, "%3d%c %16.16s %c %1.52s",
859                     (Sstack + Depth)[-i + Stacksize], VERT,
860                     (Dstack + Depth)[-i + Stacksize], VERT,
861                     yypstk( (Vstack + (Depth*Vsize)) +
862                             ((-i + Stacksize) * Vsize),
863                             (Dstack + Depth)[-i + Stacksize] ) );
864         }
865     }
866     else
867     {
868         wmove( Stack_window, Stacksize - i, 0 );
869         wclrtoeol( Stack_window );
870     }
871     wrefresh( Stack_window );
872 }
873 }
874 delay();
875 wrefresh( Stack_window );
876 Onumele = numele;
877 }
878
879 /*-----*/
880
881 PUBLIC void yy_redraw_stack()
882 {
883     /* Redraw the entire stack screen by writing out the top Stacksize elements
884     * of the stack in the stack window. Note that scrolling is turned off so
885     * that the screen won't scroll when you print the last line. Unlike
886     * yy_pstack(), this routine won't pause for a command.
887     */
888
889     int i;
890     int numele;                                /* Number of elements on the stack */
891     int *state = *P_sp;                         /* Pointer to top of state stack */
892     char **debug = *P_dsp;                      /* Pointer to top of debug stack */
893     char *value;                               /* Pointer to top of value stack */
894

```



Listing 4.11. continued...

```

895     werase ( Stack_window );
896     scrolllok ( Stack_window, FALSE );
897
898     numele = Depth - (state - Sstack);
899     value = Vstack + ((Depth - numele) * Vsize);
900
901     wmove( Stack_window, numele <= Stacksize ? Stacksize - numele : 0, 0 );
902
903     for( i=min(Stacksize, numele); --i >= 0; ++state, ++debug, value += Vsize )
904         wprintw( Stack_window, "%3d%c %16.16s %c %1.52s\n",
905                     *state, VERT,
906                     *debug, VERT, yypstk(value, *debug) );
907
908     scrolllok( Stack_window, TRUE );
909 }
```

Listing 4.12. yydebug.c— Delay and Main Control Loop

```

910 PRIVATE void delay()
911 {
912     /* Print a prompt and wait for either a carriage return or another command.
913      * Note that the time returned by time() is the elapsed time, in seconds,
914      * from 00:00:00, January 1, 1970 GMT. Since there are roughly 31,557,600
915      * seconds in a year (365.25 * 24 * 60 * 60) and the largest (signed)
916      * 32-bit long int can hold 2,147,483,647, the time won't roll over until
917      * January 18, 2038 at 2:56:02 A.M. Don't use this program on January 18,
918      * 2038 at 2:56:02 A.M.
919 */
920
921     long          start, current;
922     char          buf[80];
923     int           print_lines;
924     struct timeb time_buf;           /* defined in sys/timeb.h */
925
926     if( !Interactive )           /* n command (noninteractive) issued */
927         return;
928
929     if( !Singlestep && kbhit() )
930     {
931         /* If we're not single stepping (a 'g' command has been issued) and
932         * there's a key pressed, stop single stepping and get the character.
933         */
934         input_char();
935         Singlestep = 1;
936     }
937
938     if( !Singlestep )
939     {
940         /* If we're still doing a go command (no key was found in the previous
941         * if statement), then delay for a while. Must use two if statements
942         * here because we don't want to delay if we've just stopped go-ing.
943         * If a key is hit while we're delaying, stop looping immediately and
944         * revert back to single-step mode.
945         */
946 }
```



Listing 4.12. continued...

```

947         ftime( &time_buf );
948         start    = (time_buf.time * 1000) + time_buf.millitm;
949
950     while( 1 )
951     {
952         ftime( &time_buf );
953         current = (time_buf.time * 1000) + time_buf.millitm;
954
955         if( current - start >= Delay )
956             break;
957
958         if( kbhit() )           /* If a key is hit, stop delaying */
959         {                      /* and revert back to single-step */
960             input_char();       /* mode. */
961             Singlestep = 1;
962             break;
963         }
964     }
965     if( !Singlestep )      /* If we're still not single stepping, then */
966         return;            /* we're done (don't get a command), */
967                           /* otherwise, fall out of this block and */
968                           /* enter the command loop, below. */          */
969 }
970
971 while( 1 )
972 {
973     yyprompt( "Enter command (space to continue, ? for list): ", buf, 0 );
974
975     switch( *buf )
976     {
977         case '\0' :
978         case ' ' :
979         case '\n' :           /* singlestep */
980             goto outside;
981
982         case '?' :           /* help */
983             cmd_list();
984             NEWLINE( Prompt_window );
985             presskey();
986             yy_redraw_stack();
987             break;
988
989         case 'a':             /* abort */
990             Abort = 1;
991             Singlestep = 0;
992             goto outside;
993
994         case 'b':             /* breakpoints */
995             breakpoint();
996             yy_redraw_stack();
997             break;
998
999         case 'd':             /* set delay time */
1000
1001         if(yyprompt("Delay time (in seconds, CR=0, ESC cancels): ", buf,1))
1002             Delay = (long) (atof(buf) * 1000.0 );
1003             break;
1004

```



Listing 4.12. continued...

```

1005      case 'f':                                /* read file */
1006
1007      if( !yyprompt( "Print line numbers? (y/n, CR=y, ESC cancels): ", 
1008                           buf, 0 ) )
1009          break;
1010
1011      print_lines = *buf != 'n';
1012      if( !yyprompt( "File name or ESC to cancel: ", buf, 1 ) )
1013          break;
1014
1015      werase( Stack_window );
1016      display_file( buf, sizeof(buf), print_lines );
1017      yy_redraw_stack();
1018      break;
1019
1020      case 'g':                                /* go! */
1021          Singlestep = 0;
1022          goto outside;
1023
1024      case 'i':
1025          if( yyprompt( "Input file name or ESC to cancel: ", buf, 1 ) )
1026              new_input_file( buf );
1027          break;
1028
1029      case 'l':                                /* enable logging */
1030          to_log( buf );
1031          break;
1032
1033      case 'N':                                /* noninteractive w/o logging */
1034          Log           = NULL;
1035          No_stack_pix = 1;
1036          Interactive   = 0;
1037          Singlestep    = 0;
1038          Delay         = 0L;
1039          werase( Stack_window );
1040          goto outside;
1041
1042      case 'n':                                /* noninteractive mode w/ log */
1043          if( !Log && !to_log(buf) )
1044              break;
1045          Interactive = 0;
1046          Singlestep  = 0;
1047          Delay       = 0L;
1048          werase( Stack_window );
1049          goto outside;
1050
1051      case 'q':                                /* exit to operating system */
1052          raise( SIGINT );                      /* as if Ctrl-C was entered */
1053          exit(0);
1054
1055      case 'r':                                /* redraw the stack window */
1056          yy_redraw_stack();
1057          break;
1058
1059      case 'w':                                /* write screen to file */
1060          if( yyprompt( "Output file name or ESC to cancel: ", buf, 1 ) )
1061              write_screen( buf );
1062          break;
1063

```



Listing 4.12. continued...

```

1064     case 'x':                                /* show lexemes */
1065         yycomment( "current  [%0.*s]\n", yylen,      yytext      );
1066         yycomment( "previous [%0.*s]\n", ii_plength(), ii_ptext() );
1067         break;
1068
1069     case 0x01: yyhook_a(); break;    /* Ctrl-A debugger hook (see text) */
1070     case 0x02: yyhook_b(); break;    /* Ctrl-B */
1071
1072     default:
1073         yyprompt( "Illegal command, press any key to continue", buf, 0 );
1074         break;
1075     }
1076 }
1077 outside:
1078     werase ( Prompt_window );
1079     wrefresh( Prompt_window );
1080 }
1081
1082 /*-----*/
1083
1084 PRIVATE void cmd_list()
1085 {
1086     /* Print a list of commands in the stack window & prompt for an action. */
1087
1088     werase (Stack_window );
1089     wmove (Stack_window, 0, 0 );
1090     wprintw (Stack_window, "a (a)bort parse by reading EOI      \n");
1091     wprintw (Stack_window, "b modify or examine (b)eakpoint   \n");
1092     wprintw (Stack_window, "d set (d)elay time for go mode   \n");
1093     wprintw (Stack_window, "f read (f)ile                  \n");
1094     wprintw (Stack_window, "g (g)o (any key stops parse) \n");
1095     wprintw (Stack_window, "i change (i)nput file    \n");
1096     wmove (Stack_window, 0, 39 );
1097     wprintw (Stack_window, "l (l)o g output to file" );
1098     wmove (Stack_window, 1, 39 );
1099     wprintw (Stack_window, "n (n)o ninteractive mode" );
1100     wmove (Stack_window, 2, 39 );
1101     wprintw (Stack_window, "q (q)uit (exit to dos)" );
1102     wmove (Stack_window, 3, 39 );
1103     wprintw (Stack_window, "r (r)efresh stack window" );
1104     wmove (Stack_window, 4, 39 );
1105     wprintw (Stack_window, "w (w)rite screen to file or device\n" );
1106     wmove (Stack_window, 5, 39 );
1107     wprintw (Stack_window, "x Show current and prev. le(X)eme\n" );
1108     wmove (Stack_window, 7, (78-29)/2 );
1109     wprintw (Stack_window, "Space or Enter to single step" );
1110     wrefresh (Stack_window );
1111 }
```

One point of note is the Ctrl-A and Ctrl-B commands, processed on lines 1069 and 1070 of Listing 4.12. These commands call the yyhook_a() and yyhook_b() subroutines (in Listings 4.13 and 4.14), which do nothing at all. Their purpose is twofold. First, if you are running the parser under your compiler's debugger, these commands give you a hook into that debugger. You can set a breakpoint at yyhook_a(), and then issue a Ctrl-A to transfer from the running parser to the debugger itself. These hooks also let you add commands to the parser without having to recompile it. Since yyhook_a() and yyhook_b() are alone in separate files, customized versions that you supply are

Debugger hooks for Ctrl-A, Ctrl-B: yyhook_a(), yyhook_b().

linked rather than the default ones in the library. You can effectively add commands to the parser's debugging environment by providing your own version of one or both of these routines. For example, it's convenient to print the symbol table at various states of the compilation process so you can watch symbols being added. You can add this capability to the parser by writing a routine that prints the symbol table and calling that routine `yyhook_a()`. Thereafter, you can print the symbol table by issuing a Ctrl-A at the parser's *command* prompt.

Listing 4.13. `yyhook_a.c`—Debugger Hook 1

```
1 void yyhook_a() {} /* entered with a ^A command */
```

Listing 4.14. `yyhook_b.c`—Debugger Hook 2

```
1 void yyhook_b() {} /* entered with a ^B command */
```

Production breakpoints.

Token input:
`yy_nextoken()` calls
`yylex()`.

Listing 4.15 contains support for the input and production breakpoints. (Stack breakpoints were handled in `yy_pstack()`—the code of interest starts on line 785 of Listing 4.10, page 258). The `yy_nextoken()` function (starting on line 1112 of Listing 4.15) gets an input token from `yylex()`, though if `Abort` is true, then an *a* command has been issued, and the end-of-input marker (0) is used rather than the next token. The routine then echoes the new token to the input window, and triggers a breakpoint (on line 1142) if necessary, by setting `Singlestep` true. `Singlestep` causes `delay()` to wait for a command the next time it's called (after the next stack-window update). `yy_break()` (on line 1167) does the same thing, but for production breakpoints. It is called from the **LLama**-generated parser just before every production is applied (when a nonterminal on the stack is replaced by its right-hand side). In `occs`, it's called just before a reduction takes place. `breakpoint()` (on line 1187 of Listing 4.15) takes care of setting the breakpoints, and so forth. It processes the *b* command.

The remainder of the file (in Listing 4.16) comprises little support functions which are adequately commented and don't require additional discussion here.

Listing 4.15. `yydebug.c`—Breakpoint Support

```
1112 PUBLIC int yy_nextoken()
1113 {
1114     /* Input a token from yylex() and echo to both the token and comment
        * windows. Yy_comment() writes to the log file too. Break if the input
        * breakpoint is set and the required token has just been read. The token
        * name is centered in the token window if it's short enough. It's
        * truncated at TOKEN_WIDTH characters otherwise.
1115 */
1116     static int tok = -1; /* current token */
1117     int width;
1118     char *str;
1119     char *lexeme;
1120     char buf[ TOKEN_WIDTH ];
1121     extern char *Yy_stok[]; /* Generated by occs and llama */
1122
1123     if( tok >= 0 && (Interactive || Log) )
1124         yycomment( "Advance past %s\n", Yy_stok[tok] );
```

Listing 4.15. continued...

```

1131     lexeme = ((tok == Abort ? 0 : yylex()) == 0) ? "" : yytext;
1132
1133     if( Interactive || Log )
1134         yycomment( "Read %s <%s>\n", str = Yy_stok[tok], lexeme );
1135
1136     if( Interactive )
1137     {
1138         NEWLINE( Token_window );
1139         concat( TOKEN_WIDTH, buf, str, " ", lexeme, NULL );
1140         wprintf( Token_window, "%0.*s", TOKEN_WIDTH - 2, buf );
1141
1142         if( L_breakpoint != -1 && L_breakpoint <= yylineno )
1143         {
1144             L_breakpoint = -1;
1145             Singlestep = 1;
1146             yy_pstack( 0, 1 );
1147         }
1148         else if( (*I_breakpoint &&
1149             ( (isdigit(*I_breakpoint) && tok == atoi(I_breakpoint)) ||
1150                 !strcmp( lexeme, I_breakpoint ) ||
1151                 !strcmp( Yy_stok[tok], I_breakpoint ) )
1152             )
1153             )
1154         {
1155             Singlestep = 1;
1156             yy_pstack( 0, 1 );
1157         }
1158     }
1159 }
1160
1161     delay();
1162     return tok;
1163 }
1164
1165 /*-----*/
1166
1167 PUBLIC void yy_break( production_number )
1168 int production_number;
1169 {
1170     /* Handles production-number breakpoints. If a break is required, start
1171      * single stepping and print the stack. Stack breakpoints are handled in
1172      * yy_pstack and input breakpoints are done in yy_nextoken().
1173      *
1174      * If production_number == -1, a break is forced regardless of the value
1175      * of P_breakpoint;
1176      */
1177
1178     if( production_number == P_breakpoint || production_number == -1 )
1179     {
1180         Singlestep = 1;
1181         yy_pstack( 0, 1 );
1182     }
1183 }
1184
1185 /*-----*/
1186
1187 PRIVATE int breakpoint()
1188 {

```



Listing 4.15. continued...

```

1189     /* Set up a breakpoint by prompting the user for any required information.
1190      * Return true if we have to redraw the stack window because a help screen
1191      * was printed there.
1192      */
1193
1194     int      type;
1195     char    **p;
1196     char    buf[80];
1197     int      rval   = 0;
1198     static char *text[] =
1199     {
1200         "Select a breakpoint type (i,l,p,or s) or command (c or l):",
1201         "Type: Description:           Enter breakpoint as follows:",
1202         " i   input.....number for token value",
1203         "       or string for lexeme or token name",
1204         " l   input line read.....line number",
1205         " p   reduce by production.....number for production number",
1206         " s   top-of-stack symbol.....number for state-stack item",
1207         "       or string for symbol-stack item",
1208         " c = clear all breakpoints",
1209         " d = display (list) all breakpoints",
1210         NULL
1211     };
1212
1213     if( !yyprompt("Enter type or command, ? for help, ESC aborts: ", buf,0) )
1214         return 1;
1215
1216     if( *buf == '?' )
1217     {
1218         rval = 1;
1219         werase (Stack_window);
1220         wmove   (Stack_window, 0, 0);
1221
1222         for( p = text; *p; p )
1223             wprintw(Stack_window, "%s\n", *p++ );
1224
1225         wrefresh (Stack_window);
1226         if( !yyprompt("Enter breakpoint type or command, ESC aborts: ", buf,0))
1227             return rval;
1228     }
1229
1230     if( (type = *buf) == 'p' )
1231     {
1232         if( yyprompt("Production number or ESC to cancel: ", buf, 1 ))
1233         {
1234             if( !isdigit( *buf ) )
1235                 yyprompt("Must be a number, press any key to continue.", buf,0);
1236             else
1237                 P_breakpoint = atoi( buf );
1238         }
1239     }
1240     else if( type == 'l' )
1241     {
1242         if( yyprompt("Input line number or ESC to cancel: ", buf, 1 ))
1243             L_breakpoint = atoi( buf );
1244     }
1245     else if( type == 'i' || type == 's' )
1246     {

```



Listing 4.15. continued...

```

1247     if( yyprompt("Symbol value or ESC to cancel: ", buf, 1 ))
1248         strncpy( type == 'i' ? I_breakpoint : S_breakpoint, buf, BRKLEN );
1249     }
1250 else
1251 {
1252     switch( type )
1253     {
1254     case 'c':
1255         P_breakpoint = -1;
1256         L_breakpoint = -1;
1257         *S_breakpoint = 0;
1258         *I_breakpoint = 0;
1259         break;
1260
1261     case 'd':
1262         rval = 1;
1263         werase (Stack_window);
1264         wmove (Stack_window, 0, 0);
1265
1266         wprintw(Stack_window,
1267                 P_breakpoint == -1 ? "Production = none\n"
1268                               : "Production = %d\n", P_breakpoint);
1269
1270         wprintw(Stack_window, "Stack      = %s\n",
1271                 *S_breakpoint ? S_breakpoint : "none" );
1272
1273         wprintw(Stack_window, "Input      = %s\n",
1274                 *I_breakpoint ? I_breakpoint : "none" );
1275         wprintw(Stack_window,
1276                 I_breakpoint == 0 ? "Input line = none\n"
1277                               : "Input line = %d\n", I_breakpoint);
1278         wrefresh(Stack_window);
1279         NEWLINE (Prompt_window);
1280         presskey();
1281         break;
1282
1283     default:
1284         yyprompt("Illegal command or type, Press any key.", buf, 0);
1285         break;
1286     }
1287 }
1288
1289 return rval;
1290 }
```

Listing 4.16. yydebug.c—Input Routines

```

1291 PRIVATE int    new_input_file( buf )
1292 char   *buf;
1293 {
1294     /* Open up a new input file. Input must come from a file because the
1295     * keyboard is used to get commands. In theory, you can use both standard
1296     * input and the keyboard (by opening the console as another stream), but I
1297     * had difficulties doing this in a portable way, and eventually gave up.
1298     * It's not that big a deal to require that test input be in a file.
1299     */
1300 }
```

Listing 4.16. continued...

```

1301     NEWLINE( Prompt_window );
1302     wrefresh( Prompt_window );
1303
1304     if( ii_newfile( buf ) != -1 )
1305         Inp_fm_file = 1;
1306     else
1307     {
1308         wprintw(Prompt_window, "Can't open %s.", buf );
1309         presskey();
1310     }
1311     return Inp_fm_file;
1312 }
1313
1314 /*-----*/
1315
1316 PRIVATE FILE *to_log( buf )
1317 char *buf;
1318 {
1319     /* Set up everything to log output to a file (open the log file, etc.). */
1320
1321     if( !yyprompt("Log-file name (CR for \"log\\\", ESC cancels): ", buf,1) )
1322         return NULL;
1323
1324     if( !*buf )
1325         strcpy( buf, "log" );
1326
1327     if( !(Log = fopen( buf, "w")) )
1328     {
1329         NEWLINE(Prompt_window );
1330         wprintw(Prompt_window, "Can't open %s", buf );
1331         presskey();
1332         return NULL;
1333     }
1334
1335     if( !yyprompt("Log comment-window output? (y/n, CR=y): ", buf,0) )
1336         return NULL;
1337     else
1338         No_comment_pix = (*buf == 'n');
1339
1340     if( !yyprompt( "Print stack pictures in log file? (y/n, CR=y): ",buf,0) )
1341         return NULL;
1342
1343     if( !(No_stack_pix = (*buf == 'n')) )
1344     {
1345         if( !yyprompt( "Print stacks horizontally? (y/n, CR=y): ",buf,0) )
1346             return NULL;
1347
1348         if( Horiz_stack_pix = (*buf != 'n') )
1349         {
1350             if( !yyprompt("Print SYMBOL stack (y/n, CR=y): ",buf,0) )
1351                 return NULL;
1352             Sym_pix = (*buf != 'n');
1353
1354             if( !yyprompt("Print PARSE stack (y/n, CR=y): ",buf,0) )
1355                 return NULL;
1356             Parse_pix = (*buf != 'n');
1357
1358             if( !yyprompt("Print VALUE stack (y/n, CR=y): ",buf,0) )
1359                 return NULL;
1360             Attr_pix = (*buf != 'n');

```



Listing 4.16. continued...

```

1361         }
1362     }
1363     return Log;
1364 }
1365
1366 /*-----*/
1367
1368 PRIVATE int input_char( void )
1369 {
1370     /* Get a character from the input window and echo it explicitly. If we've
1371      * compiled under Unix, reset the character-available flag.
1372      */
1373
1374     int c;
1375
1376     if( (c = wgetch(Prompt_window) & 0x7f) != ESC )
1377         waddch( Prompt_window, c );
1378
1379     UX( Char_avail = 0; )
1380     return c;
1381 }
1382
1383 /*-----*/
1384
1385 PUBLIC int yyprompt( prompt, buf, getstring )
1386 char *prompt, *buf;
1387 int getstring;    /* get entire string (as compared to a single character */
1388 {
1389     /* Print a prompt and then wait for a reply, load the typed characters into
1390      * buf. ^H (destructive backspace) is supported. An ESC causes yyprompt()
1391      * to return 0 immediately, otherwise 1 is returned. The ESC is not put
1392      * into the buffer. If "getstring" is true, an entire string is fetched and
1393      * prompt returns when the newline is typed. If "getstring" is false, then
1394      * one character is fetched and yyprompt() returns immediately after getting
1395      * that character. Leading and trailing white space is stripped from the
1396      * line if getstring is true. (You can get a single space character if it's
1397      * false, however.)
1398      */
1399
1400     register int c;
1401     int y, x;
1402     char *startbuf = buf;
1403
1404     NEWLINE( Prompt_window );
1405     wprintw( Prompt_window, "%s", prompt );
1406     wrefresh( Prompt_window );
1407
1408     if( !getstring )
1409         c = *buf++ = input_char();
1410     else
1411     {
1412         while( (c = input_char()) != '\n' && c != ESC )
1413         {
1414             if( isspace(c) && buf==startbuf )
1415                 continue;                                /* skip leading white space */
1416
1417             if( c != '\b' )
1418                 *buf++ = c;

```



Listing 4.16. continued...

```

1419         else                                /* handle destructive backspace */
1420     {
1421         getyx( Prompt_window, y, x );
1422
1423         if( buf <= startbuf )
1424             wmove ( Prompt_window, y, x+1 );
1425         else
1426         {
1427             waddch( Prompt_window, ' ' );
1428             wmove ( Prompt_window, y, x );
1429             --buf;
1430         }
1431     }
1432     wrefresh( Prompt_window );
1433 }
1434
1435     while( isspace( buf[-1] )  &&  buf > startbuf )
1436     --buf;                            /* Strip trailing white space */
1437 }
1438 *buf = 0;
1439 return (c != ESC);
1440 }
1441
1442 /*-----*/
1443
1444 PRIVATE void presskey()
1445 {
1446     /* Ask for a key to be pressed and wait for it. Note that this command
1447      * does a refresh, but it intentionally does not clear the window before
1448      * printing the prompt.
1449      */
1450
1451     wprintw ( Prompt_window, " Press any key: " );
1452     wrefresh( Prompt_window );
1453     input_char();
1454 }
```

4.10 LLama—Implementing an LL(1) Parser-Generator

The remainder of this chapter discusses how **LLama** works. Much of the following code is used by both **LLama** and **occ**s (all the files whose names start with *ll* are used only by **LLama**, everything else is used jointly). As usual, I expect you to actually read the code (I've not repeated in the text those things that are adequately described in comments in the code itself). I also expect you to be familiar with the set and hash-table functions described in Appendix A.

4.10.1 LLama's Parser

LLama's own, recursive-descent parser.

This section describes **LLama**'s own parser. It provides a good example of a working recursive-descent compiler for a small programming language (**LLama** is itself a compiler, after all). It translates a very high-level language (a grammatical description of a programming language) to a high-level language (C).

I built **LLama** in a two-step process, initially developing yet another recursive-descent parser (the last one we'll look at in this book) that supports a bare-bones input language. I then used this recursive-descent version of **LLama** to rebuild its own parser, substituting the **LLama** output file for the recursive-descent parser when I linked the new version. This process is typical of how a language is brought up on a new machine. You first construct a simple compiler for a subset of the target language using the tools at hand—typically an assembler. The output of this compiler is typically the same assembly language that the compiler itself is written in. Using your subset language, you then write a compiler for a more complete version of the language.

Bootstrapping a compiler using language subsets.

For example, you could start writing a C compiler with a very small subset of C, written in assembly language—an adequate subset would support simple expressions (no fancy operators like conditionals), global variables of type **int** and **char** (but no local variables), one-dimensional arrays of **int** and **char**, simple control-flow statements (**while**, **if**, and **else** are adequate), and some sort of block structure. Note that, at this level, a compiler is little more than a macro translator. (See [Angermeyer], pp. 51–92, where all of the foregoing are implemented as macros using the Microsoft macro assembler, MASM.)

In the next step, you write a larger subset of C using the language defined in the previous step, and compile it using the subset-of-C compiler. Typically, this second level would add better typing (introducing structures, for example), support local variables and subroutine arguments, and so forth. Continuing in this manner, you can bootstrap yourself up to a full language implementation.

LLama, like most compilers, is easily divided into several distinct phases. The lexical analyzer is created with L^EX. (The specification is in Listing 4.17 and the required token definitions are in Listing 4.18.) Note that white space (a comment is treated like white space) is ignored if the global variable **Ignore** (declared on line 27) is true, otherwise a **WHITE SPACE** token is returned. The variable is modified using the access routines on lines 215 and 216. Also note how actions are processed on lines 73 to 153. The entire action is treated as a single token. L^EX recognizes an initial open brace, and the associated code absorbs everything up to and including the matching close brace. That is, the code doesn't terminate until a close brace is found at nesting level 1, as defined by **nestlev**. Other close braces are absorbed into the lexeme.

LLama's lexical analyzer, tokens, L^EX input specification.

Listing 4.17. *parser.lex*—L^EX Input File for **occ**/**LLama** Lexical Analyzer

```

1  %
2  #include <stdio.h>
3  #include <tools/debug.h>
4  #include <tools/hash.h>
5  #include <tools/set.h>
6  #include "llout.h"
7
8  #define CREATING_LLAMA_PARSER /* Suppress various definitions in parser.h */
9  #include "parser.h"           /* that conflict with LEX-generated defs. */
10
11 /* -----
12 * Lexical analyzer for both llama and yacc. Note that llama doesn't support
13 * %left, %right, %noassoc, or %prec. They are recognized here so that we can
14 * print an error message when they're encountered. By the same token, yacc
15 * ignores the %synch directive. Though all legal llama input files can be
16 * processed by yacc, the reverse is not true.
17 *
18 * -----
19 * Whitespace, comments, and otherwise illegal characters must be handled
20 * specially. When we're processing code blocks, we need to get at

```



Listing 4.17. continued...

```

21   * the characters so that they can be passed to the output, otherwise, the
22   * characters should be ignored. The ws() and nows() subroutines (at the
23   * bottom of the file) switch between these behaviors by changing the value
24   * if Ignore. If Ignore is true, white space is ignored.
25   */
26
27 PRIVATE int    Ignore = 0;
28 PRIVATE int    Start_line;      /* starting line number */
29
30 /* -----
31  * Prototypes for functions at the bottom of this file:
32 */
33
34 void stripcr P(( char *src )); /* Remove carriage returns (but not      */
35                                /* linefeeds) from src.                  */
36 void nows     P(( void ));    /* Ignore white space, etc               */
37 void ws       P(( void ));    /* Don't ignore white space, etc        */
38 %}
39
40 c_name  [A-Za-z_] [A-Za-z_0-9]*
41 %%
42
43 "/*"           { /* Absorb a comment (treat it as WHITESPACE) */
44
45     int i;
46     int start = yylineno;
47
48     while( i = input() )
49     {
50         if( i < 0 )
51         {
52             ii_unterm();
53             ii_flush(1);
54             ii_term();
55             lerror(NONFATAL, "Comment starting on line %d " \
56                   "too long, truncating\n", start);
57         }
58         else if( i == '*' && ii_lookahead(1) == '/' )
59         {
60             input();
61             stripcr( yytext );
62
63             if( Ignore ) goto end;
64             else         return WHITESPACE;
65         }
66     }
67
68     lerror(FATAL, "End of file encountered in comment\n");
69
70     end:;
71 }
72
73 /* Suck up an entire action. Handle nested braces here.
74  * This code won't work if the action is longer than the
75  * buffer length used by the input functions. If this is a
76  * problem, you have to allocate your own buffer and copy
77  * the lexeme there as it's processed (in a manner similar to
78  * the ${ processing, below). If space is really a problem,
79  * the code blocks can be copied to a temporary file and the
80  * offset to the start of the text (as returned by ftell()) */

```



Listing 4.17. continued...

```

81                                * can be stored rather than the string itself.
82                                */
83
84    \{
85        int i;
86        int nestlev;           /* brace-nesting level      */
87        int lbl;               /* previous character      */
88        int lb2;               /* character before that   */
89        int in_string;         /* processing string constant */
90        int in_char_const;     /* processing char. constant */
91        int in_comment;        /* processing a comment     */
92
93        lbl = lb2 = 0;
94        in_string = 0;
95        in_char_const = 0;
96        in_comment = 0;
97        Start_line = yylineno;
98
99        for( nestlev=1, i=input(); lb2=lb1, lbl=i )
100    {
101            if( lb2=='\n' && lbl=='%' && i=='%' )
102                lerror( FATAL,
103                         "%%% in code block starting on line %d\n",
104                         Start_line );
105
106            if( i < 0 )           /* input-buffer overflow */
107            {
108                ii_unterm();
109                ii_flush(1);
110                ii_term();
111                lerror( FATAL,
112                         "Code block starting on line %d too long.\n",
113                         Start_line );
114            }
115
116            /* Take care of \{, "{", '(', '\}, "}"', ')' */
117
118            if( i == '\\\' )
119            {
120                if( !(i = input()) ) /* discard backslash */
121                    break;
122                else
123                    continue;          /* and following char */
124            }
125
126            if( i == '"' && !(in_char_const || in_comment) )
127                in_string = !in_string;
128
129            else if( i == '\'\' && !(in_string || in_comment) )
130                in_char_const = !in_char_const;
131
132            else if( lbl == '/' && i == '*' && !in_string )
133                in_comment = 1;
134
135            else if( lbl == '*' && i == '/' && in_comment )
136                in_comment = 0;
137

```



Listing 4.17. continued...

```

138                     if( !(in_string || in_char_const || in_comment) )
139                     {
140                         if( i == '{' )
141                             ++nestlev;
142
143                         if( i == '}' && --nestlev <= 0 )
144                         {
145                             stripchr( yytext );
146                             return ACTION;
147                         }
148                     }
149
150                     lerror(FATAL, "EOF in code block starting on line %d\n",
151                                         Start_line );
152
153                 }
154
155             ^"%%"
156             return SEPARATOR;           /* Must be anchored because */
157                                         /* it can appear in a printf */
158                                         /* statement. */           */
159             "%{["[\s\t]*"
160             {
161                 /* Copy a code block to the output file. */
162
163                 int c, looking_for_brace = 0;
164
165                 #undef output          /* replace macro with function */
166                                         /* in main.c */           */
167                 if( !No_lines )
168                     output( "\n#line %d \"%s\"\n",
169                                         yylineno, Input_file_name );
170
171                 while( c = input() )      /* while not at end of file */
172                 {
173                     if( c == -1 )        /* buffer is full, flush it */
174                         ii_flushbuf();
175
176                     else if( c != '\r' )
177                     {
178                         if( looking_for_brace ) /* last char was a % */
179                         {
180                             if( c == '}' ) break;
181                             else           output( "%%c", c );
182                         }
183                         else
184                         {
185                             if( c == '%' ) looking_for_brace = 1;
186                             else           output( "%c", c );
187                         }
188                     }
189                 }
190
191             <{c_name}>
192             return FIELD;           /* for occs only */
193             "%union"
194             "%token"
195             "%term"
196             "%type"
197             "%synch"
198             "%left"
199             return PERCENT_UNION;   /* for occs only */
200             |
201             return TERM_SPEC;
202             return TYPE;           /* for occs only */
203             return SYNCH;          /* for llama only */
204             return LEFT;           /* for occs only */
205
206             }
207
208         }
209
210     }
211
212     return CODE_BLOCK;
213
214 }
```



Listing 4.17. continued...

```

198 "%right"           return RIGHT;          /* for occs only */      */
199 "%nonassoc"         return NONASSOC;       /* for occs only */     */
200 "%prec"             return PREC;          /* for occs only */      */
201 "%start"            return START;          /* for error messages only */
202 ":"                return COLON;          ;
203 "|"                return OR;              ;
204 ";"                return SEMI;           ;
205 "["                return START_OPT;      ;
206 "]"                |
207 "]**"              return END_OPT;
208
209 [^\x00-\s%\{\}[\]]()*:|;,<>]+    return NAME;
210 \x0d                  ; /* discard carriage return (\r) */
211 [\x00-\x0c\x0e-\s]        if( !Ignore ) return WHITESPACE;
212 %%
213 /*-----*/
214
215 PUBLIC void nows() { Ignore = 1; }      /* Ignore white space, etc. */
216 PUBLIC void ws () { Ignore = 0; }        /* Don't ignore white space, etc. */
217
218 PUBLIC int start_action() /* Return starting line number of most */
219 {                           /* recently read ACTION block */
220     return Start_line;
221 }
222 /*-----*/
223
224 PRIVATE void stripcr( src ) /* Remove all \r's (but not \n's) from src. */
225 char *src;
226 {
227     char *dest;
228     for( dest = src ; *src ; src++ )
229         if( *src != '\r' )
230             *dest++ = *src;
231     *dest = '\0';
232 }
```

LLama's parser uses the L^EX-generated lexical analyzer to parse the LLama input file, creating a symbol table with what it finds there. A *symbol table* is a large data structure, indexed by symbol name (the lexeme associated with the name in the input file). A good analogy is an array of structures, one field of which is a string holding the symbol name. The other fields in the structure hold things like the numeric value assigned to each symbol (the token value) and the symbol's type (in the current case, the type is *terminal*, *nonterminal*, or *action*, as compared to the more usual `int`, `long`, and so forth). LLama's symbol also stores tokenized versions of the right-hand sides as elements of the symbol-table entry for each nonterminal. I'll discuss the details of this process in a moment. The LLama code-generation phase is passed the symbol table, which in this case serves as a representation of the entire input language, and outputs code—the parse tables needed to parse an input file in the target language. It also copies the output parser from the template file to the output file.

The `parser.h` file, which starts in Listing 4.19, is `#included` in most of the files that comprise both LLama and occs. It starts out on lines six to 12 with conditional-compilation macros that work like the `D()` macro in `<tools/debug.h>` (described in Appendix A). The argument to the `LL()` macro compiles only if you're making LLama (I usually define `LLAMA` on the compiler's command line with a `cc -DLLAMA`), the opposite applies to the `OX()` macro. The listing also contains a few predefined exit codes.

LLama's symbol-table.

Conditional compilation
macros: `LL()`, `OX()`.

Listing 4.18. *llout.h*—LLama Token Definitions

```

1 #define _EOI_          0
2 #define ACTION         1
3 #define CODE_BLOCK     2
4 #define COLON          3
5 #define END_OPT        4
6 #define FIELD          5
7 #define LEFT           6
8 #define NAME           7
9 #define NONASSOC       8
10 #define OR             9
11 #define OTHER          10
12 #define PREC           11
13 #define RIGHT          12
14 #define SEMI           13
15 #define SEPARATOR      14
16 #define START          15
17 #define START_OPT      16
18 #define SYNCH          17
19 #define TERM_SPEC      18
20 #define TYPE           19
21 #define PERCENT_UNION 20
22 #define WHITESPACE     21

```

Listing 4.19. *parser.h*—Compilation Directives and Exit Stati

```

1 /*      PARSER.H      This file contains those #defines, etc., that
2 *                                are used by both llama and yacc. There's also
3 *                                a yacc.h file that's used only by the yacc code.
4 */
5
6 #ifdef LLAMA
7 #define LL(x) x
8 #define OX(x)
9 #else
10 #define LL(x)
11 #define OX(x) x
12 #endif
13
14 /*****
15 * Various error exit stati. Note that other errors found while parsing cause
16 * llama to exit with a status equal to the number of errors (or zero if
17 * there are no errors).
18 */
19
20 #define EXIT_ILLEGAL_ARG      255    /* Illegal command-line switch      */
21 #define EXIT_TOO_MANY          254    /* Too many command-line args      */
22 #define EXIT_NO_DRIVER         253    /* Can't find llama.par          */
23 #define EXIT_OTHER              252    /* Other error (syntax error, etc.) */
24 #define EXIT_USR_ABRT          251    /* Ctrl-Break                      */

```

Token values: MINNON-
TERM, MINACT.

Parser.h continues in Listing 4.20 with the definitions of the numbers that represent terminals, nonterminals, and actions. You can change the values of MINNONTERM and MINACT (on lines 29 and 30), but they must always be in the same relative order. The

following must hold:

$$(0 = \text{EOI}) < (\text{MINTERM} = 1) < \text{MINNONTERM} < \text{MINACT}$$

Zero is reserved for the end-of-input marker (`_EOI_`), and `MINTERM` must be 1. Also note that there must be at least one hole between the maximum terminal and minimum nonterminal values. This hole is generated with the -2 in the `MAXNONTERM` definition on line 41. It's required because the symbolic value used for ϵ is one more than the largest number actually used to represent a terminal symbol; the hole guarantees that there's enough space for the ϵ . (`EPSILON` is defined on line 65.) The listing finishes with a few miscellaneous definitions for various file names, and so forth.

Listing 4.20. `parser.h`— Numeric Limits for Token Values

```

25 #define MAXNAME      32 /* Maximum length of a terminal or nonterminal name */
26 #define MAXPROD      512 /* Maximum number of productions in the input grammar */
27
28 #define MINTERM      1 /* Token values assigned to terminals start here */
29 #define MINNONTERM  256 /* nonterminals start here */
30 #define MINACT       512 /* acts start here */
31
32                                     /* Maximum numeric values used for terminals
33                                     * and nonterminals (MAXTERM and MINTERM), as
34                                     * well as the maximum number of terminals and
35                                     * nonterminals (NUMTERMS and NUMNONTERMS).
36                                     * Finally, USED_TERMS and USED_NONTERMS are
37                                     * the number of these actually in use (i.e.
38                                     * were declared in the input file).
39                                     */
40
41 #define MAXTERM      (MINNONTERM -2)
42 #define MAXNONTERM   (MINACT    -1)
43
44 #define NUMTERMS     ((MAXTERM-MINTERM) +1)
45 #define NUMNONTERMS  ((MAXNONTERM-MINNONTERM)+1)
46
47 #define USED_TERMS   ((Cur_term      - MINTERM)      +1)
48 #define USED_NONTERMS ((Cur_nonterm - MINNONTERM) +1)
49
50                                     /* These macros evaluate to true if x represents
51                                     * a terminal (ISTERM), nonterminal (ISNONTERM)
52                                     * or action (ISACT)
53                                     */
54
55 #define ISTERM(x)     ((x) && (MINTERM      <= (x)->val && (x)->val <= MAXTERM      ))
56 #define ISNONTERM(x)   ((x) && (MINNONTERM <= (x)->val && (x)->val <= MAXNONTERM))
57 #define ISACT(x)       ((x) && (MINACT      <= (x)->val
58
59                                     /* Epsilon's value is one more than the largest
60                                     * terminal actually used. We can get away with
61                                     * this only because EPSILON is not used until
62                                     * after all the terminals have been entered
63                                     * into the symbol table.
64                                     */
65 #define EPSILON   (Cur_term+1)
66
67                                     /* The following macros are used to adjust the
68                                     * nonterminal values so that the smallest
69                                     * nonterminal is zero. (You need to do this
70                                     * when you output the tables). ADJ_VAL does
71                                     * the adjustment, UNADJ_VAL translates the
72                                     */

```



Listing 4.20. continued...

```

71                                     * adjust value back to the original value.
72                                     */
73
74 #define ADJ_VAL(x)    ((x)-MINNONTERM )
75 #define UNADJ_VAL(x)  ((x)+MINNONTERM )
76
77 /*-----*/
78
79 #define NONFATAL      0      /* Values to pass to error() and lerror(), */
80 #define FATAL         1      /* defined in main.c. */
81 #define WARNING       2
82
83 #define DOLLAR_DOLLAR ((unsigned)^0 >>1)      /* Passed to do_dollar() to */
84                                         /* indicate that $$ is to */
85                                         /* be processed. */
86
87 /*-----*/
88
89 #ifdef LLAMA
90 #  define TOKEN_FILE   "llout.h"          /* Various file names:          */
91 #  define PARSE_FILE   "llout.c"          /* output file for token #defines */
92 #  define SYM_FILE     "llout.sym"        /* output file for parser       */
93 #  define DOC_FILE     "llout.doc"        /* output file for symbol table */
94 #  define DEF_EXT      "lma"              /* LALR(1) State machine description */
95 #  define PAR_TEMPL    "llama.par"        /* foo.lma is default input extension */
96 #  define PROG_NAME    "llama"           /* template file for parser      */
97#else
98 #  define TOKEN_FILE   "yyout.h"          /* output file for token #defines */
99 #  define PARSE_FILE   "yyout.c"          /* output file for parser       */
100 #  define ACT_FILE     "yyact.c"          /* Used for output if -a specified */
101 #  define TAB_FILE     "yyoutab.c"        /* output file for parser tables (-T) */
102 #  define SYM_FILE     "yyout.sym"        /* output file for symbol table */
103 #  define DOC_FILE     "yyout.doc"        /* LALR(1) State machine description */
104 #  define DEF_EXT      "ox"               /* foo.ox is default input extension */
105 #  define PAR_TEMPL    "occs.par"         /* template file for PARSE_FILE */
106 #  define ACT_TEMPL   "occs-act.par"     /* template file for ACT_FILE   */
107 #  define PROG_NAME    "occs"
108#endif

```

Definitions for output functions, YY_TTYPE

Symbol-table data structures.

Listing 4.21 shows several definitions used by the output functions. The **typedefs** for YY_TTYPE on lines 115 and 116 are used only to predict the size of the output transition tables. They should agree with the the YY_TTYPE definitions in the parser-template files.

The next part of *parser.h*, the SYMBOL and PRODUCTION structures declared in Listing 4.22, are used for LLama's symbol table. Symbol tables tend to be the most complex data structure in the compiler, and the LLama symbol table is no exception. Figure 4.9 shows a symbol table containing an entry for the following three productions:

$$\begin{array}{rcl}
 \textit{expr} & \rightarrow & \textit{term} \textbf{ PLUS } \textit{expr} \\
 & | & \textit{term} \\
 \textit{term} & \rightarrow & \textbf{NUMBER} \{ \textit{create_tmp}(\textit{yytext}); \}
 \end{array}$$

For clarity, I've left out unimportant fields, and NULL pointers are just left blank in the picture.

SYMBOL structure.

The symbol table itself is made up of SYMBOL structures, shown in the dashed box in Figure 4.9. This box is actually a hash table, put together with the functions described in

Listing 4.21. *parser.h*— Simple Types

```

109 /* The following are used to define types of the OUTPUT transition tables. The
110 * ifndef takes care of compiling the llama output file that will be used to
111 * recreate the llama input file. We must let the llama-generated definitions
112 * take precedence over the the default ones in parser.h in this case.
113 */
114 #ifndef CREATING_LLAMA_PARSER
115     LL( typedef unsigned char      YY_TTYPE ;          )
116     OX( typedef int            YY_TTYPE ;          )
117 #endif

```

Listing 4.22. *parser.h*— The SYMBOL and PRODUCTION Data Structures

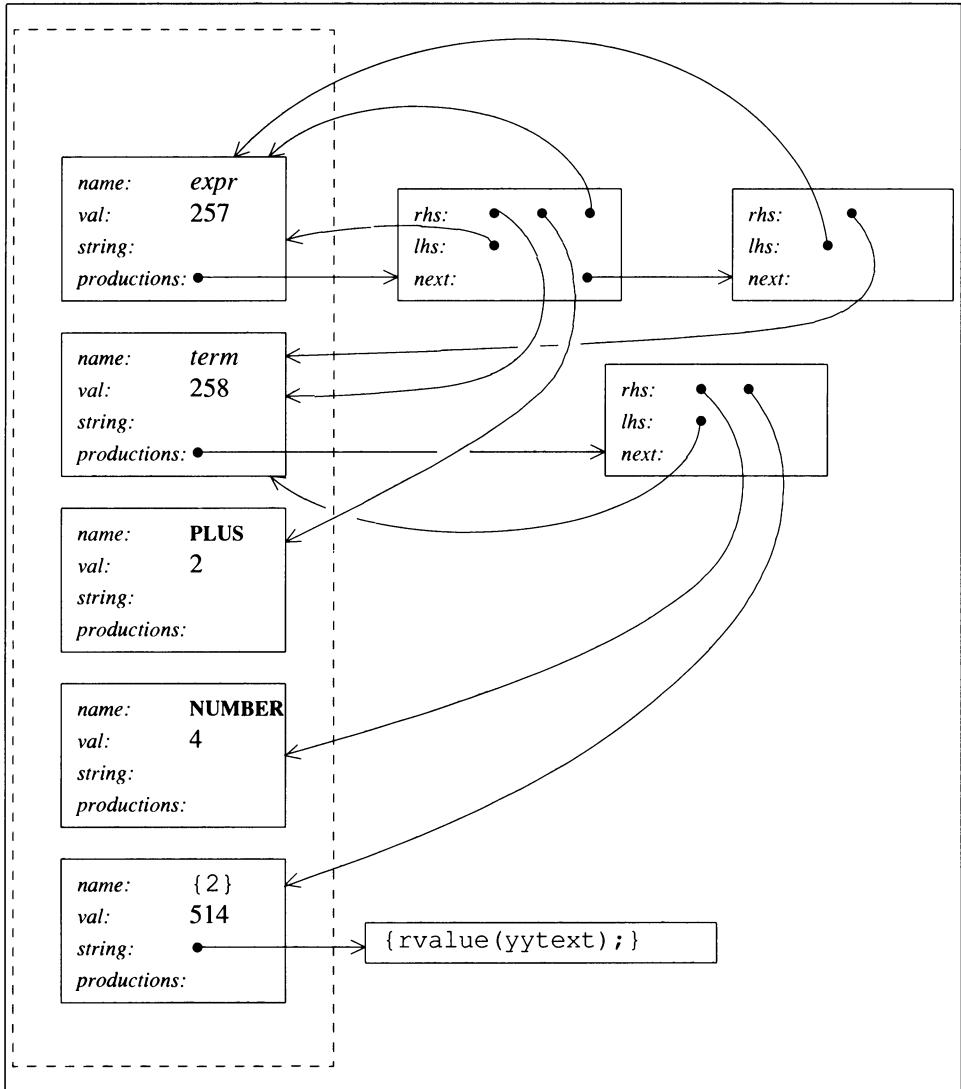
```

118 /* SYMBOL structure (used for symbol table). Note that the name itself */
119 /* is kept by the symbol-table entry maintained by the hash function. */
120
121 #define NAME_MAX 32                                /* Max name length + 1 */
122
123 typedef struct _symbol_
124 {
125     char           name [ NAME_MAX ];             /* symbol name. Must be first */
126     char           field [ NAME_MAX ];            /* %type <field> */
127     unsigned        val;                          /* numeric value of symbol */
128     unsigned        used;                         /* symbol used on an rhs */
129     unsigned        set;                          /* symbol defined */
130     unsigned        lineno;                        /* input line num. of string */
131     unsigned char  *string;                       /* code for actions. */
132     struct _prod_ *productions;                  /* right-hand sides if nonterm*/
133     SET              *first;                        /* FIRST set */
134     LL( SET          *follow; )                   /* FOLLOW set */
135 } SYMBOL;
136
137 #define NULLABLE(sym) ( ISNONTERM(sym) && MEMBER((sym)->first, EPSILON) )
138 /-----*/
139 /* PRODUCTION Structure. Represents right-hand sides. */
140
141 #define MAXRHS   31      /* Maximum number of objects on a right-hand side */
142 #define RHSBITS  5       /* Number of bits required to hold MAXRHS */
143
144 typedef struct _prod_
145 {
146     unsigned        num;                          /* production number */
147     SYMBOL         *rhs[ MAXRHS + 1];            /* Tokenized right-hand side */
148     SYMBOL         *lhs;                          /* Left-hand side */
149     unsigned char  rhs_len;                      /* # of elements in rhs[] array */
150     unsigned char  non_acts;                     /* " that are not actions */
151     SET              *select;                      /* LL(1) select set */
152     struct _prod_ *next;                        /* pointer to next production */
153                                         /* for this left-hand side. */
154     OX( int        prec; )                      /* Relative precedence */
155 }
156 } PRODUCTION;

```

Appendix A. A SYMBOL contains the name of the symbol (name) and the internal numeric value (val). If the symbol is an action, the string field points at a string that holds the code specified in the LLama input file. In the case of a nonterminal, the

Figure 4.9. Structure of the LLama Symbol Table



PRODUCTION structure.

productions field points at the head of a linked list of PRODUCTION structures, one element of the list for each right-hand side. The right-hand side is represented within the PRODUCTION as an array of pointers to symbol-table entries for each element of the right-hand side. The size of this array obviously limits the size of the right-hand side. I opted for this approach rather than keeping an array of numeric values for the symbols, because the latter would require an additional table lookup every time I had to access a symbol on a right-hand side. Note that the PRODUCTIONS form a doubly linked list in that the lhs field of each PRODUCTION points back to the production's left-hand side. The next field is just the next element in the linked list.

Other fields in the PRODUCTION are the production number (num)—an arbitrary, but unique, number assigned to each production; rhs_len is the number of symbols in the rhs[] array; non_acts is the number of elements in rhs[] that are not actions; finally, select is the LL(1) select set. (I'll discuss exactly what this is in a moment.) The prec field holds the relative precedence level of the current production if you're compiling for `occs`. This number is usually the precedence level of the rightmost

terminal on the right-hand side, but it can be changed with a %prec directive.

Other fields in the SYMBOL structure are: lineno—the input line number on which the symbol is defined; used is set true when a symbol is used on a right-hand side; and set is set true when a symbol is defined with a %term or by appearing as a left-hand side. It's a hard error for a symbol to be used without having been defined (set). It's just a warning if the symbol is defined but not used, however. Actions are set implicitly by being used. I'll discuss the first and follow fields in a moment.

The next data structure of interest is the PREC_TAB structure defined on lines 159 to 164 of Listing 4.23. This structure is used only by occs—it holds precedence and associativity information for the terminal symbols. An array of these (defined on line 200) is indexed by nonterminal value. This information could also be incorporated into the SYMBOL structure, but it's easier to keep it as a separate table. The remainder of *parser.h* comprises global-variable and subroutine declarations. Space is allocated for the variables if ALLOCATE is defined before *parser.h* is #included.

Precision table:
PREC_TAB

Listing 4.23. *parser.h*—Declarations and Type Definitions

```

157 #ifdef OCCS
158
159 typedef struct _prectab_
160 {
161     unsigned char level; /* Relative precedence, 0=none, 1=lowest      */
162     unsigned char assoc; /* associativity: 'l'=left,'r'=right,'\0'=none */
163
164 } PREC_TAB;
165
166 #define DEF_FIELD "yy_def"          /* Field name for default field in a      */
167                      /* *union.                                     */
168 #endif
169 #ifdef ALLOCATE
170 #define CLASS             /* If ALLOCATE is true, allocate space and */
171 #define I(x) x           /* activate the initializer, otherwise the   */
172 #else
173 #define CLASS extern    /* storage class is extern and the         */
174 #define I(x)             /* initializer is gobbled up.               */
175 #endif
176
177
178 CLASS int Debug           I( = 0 ); /* The following are set in      */
179 CLASS char *Input_file_name I( = "console" ); /* main.c, mostly by command- */
180 CLASS int Make_actions    I( = 1 ); /* line switches.               */
181 CLASS int Make_parser     I( = 1 ); /* */
182 CLASS int Make_yyoutab   I( = 0 ); /* */
183 CLASS int No_lines        I( = 0 ); /* */
184 CLASS int No_warnings     I( = 0 ); /* */
185 CLASS FILE *Output       I( = stdout ); /* */
186 CLASS int Public          I( = 0 ); /* */
187 CLASS int Symbols         I( = 0 ); /* */
188 CLASS int Threshold       I( = 4 ); /* */
189 CLASS int Uncompressed    I( = 0 ); /* */
190 CLASS int Use_stdout      I( = 0 ); /* */
191 CLASS int Verbose         I( = 0 ); /* */
192
193
194 CLASS SYMBOL *Terms[ MINACT ]; /* This array is indexed by terminal or */
195                                /* nonterminal value and evaluates to a */
196                                /* pointer to the equivalent symbol-table */
197                                /* entry. */
198

```



Listing 4.23. continued...

```

199
200 OX( CLASS PREC_TAB Precedence[ MINNONTTERM ]; ) /* Used only by occs. Holds */
201                                         /* relative precedence and */
202                                         /* associativity information */
203                                         /* for both terminals and */
204                                         /* nonterminals. */
205
206 LL( CLASS SET *Synch; ) /* Error-recovery synchronization */
207                                         /* set (specified with %synch). */
208                                         /* Synch is initialized in acts.c */
209                                         /* It is used only by llama. */
210 CLASS char      *Template I(=PAR_TEMPL); /* Template file for the parser; */
211                                         /* can be modified in main.c */
212 CLASS HASH_TAB *Symtab;           /* The symbol table itself */
213                                         /* (initialized in yyact.c) */
214 CLASS SYMBOL   *Goal_symbol I(=NULL); /* Pointer to symbol-table entry */
215                                         /* for the start (goal) symbol */
216                                         /* The following are used */
217                                         /* by the acts [in llact.c] */
218 CLASS int       Cur_term     I(= MINTERM-1); /* Current terminal */
219 CLASS int       Cur_nonterm  I(= MINNONTTERM-1); /* " nonterminal */
220 CLASS int       Cur_act      I(= MINACT-1); /* " action */
221 CLASS int       Num_productions I( = 0 ); /* Number of productions */
222
223 #undef CLASS
224 #undef I
225 /-----*/
226
227 #define outc(c) putc(c,Output); /* Character routine to complement */
228                                         /* output() in main.c */

```

LLama's parser does one thing: it loads the symbol table with the grammar represented by the input file. The process is analogous to creating a physical parse tree—the input language can be represented entirely by the symbol table, just as normal programming language can be represented by the parse tree.

An augmented grammar for the simple, recursive-descent version of the **LLama** parser is shown in Table 4.19. The **LLama** input file in Listing 4.24 specifies the complete input grammar for both **LLama** and **occs**. I'll build a recursive-descent **LLama** first, and then build a more complete parser later, using the stripped-down version of **LLama** for this purpose.

LLama input language:
recursive-descent grammar, LLama input file.

Recursive-descent parser
for LLama.

The recursive-descent parser for **LLama** is in Listing 4.25. It is a straightforward representation of the grammar in Table 4.19.

Table 4.19. LLama Grammar (Small Version)

Production		Implemented in this subroutine
<i>spec</i>	\rightarrow <i>definitions</i> {first_sym()} body stuff	yyparse()
<i>definitions</i>	\rightarrow TERM_SPEC <i>tname</i> <i>definitions</i> CODE_BLOCK <i>definitions</i> SYNCH <i>sname</i> <i>definitions</i> SEPARATOR EOI	definitions() definitions() definitions() definitions() definitions() definitions() definitions() definitions() definitions() definitions() definitions() definitions() body() body() body() body() body() body() body() body() body() right_sides() right_sides() rhs() rhs()
<i>sname</i>	\rightarrow NAME {add_synch(yytext)} <i>sname</i>	
<i>tname</i>	\rightarrow NAME {make_term(yytext)} <i>tname</i>	
<i>body</i>	\rightarrow <i>rule</i> <i>body</i> <i>rule</i> SEPARATOR <i>rule</i> EOI	body() body() body() body()
<i>rule</i>	\rightarrow NAME {new_nonterm(yytext,1)} COLON <i>right_sides</i> E	body() body() body()
<i>right_sides</i>	\rightarrow {new_rhs()} <i>rhs</i> OR <i>right_sides</i> {new_rhs()} <i>rhs</i> SEMI	right_sides() right_sides()
<i>rhs</i>	\rightarrow NAME {add_to_rhs(yytext,0)} <i>rhs</i> ACTION {add_to_rhs(yytext,start_action())} <i>rhs</i>	rhs() rhs()

Listing 4.24. *parser.lma*— LLama Input Specification for Itself

```

1  %{
2  #include <stdarg.h>
3  #include <tools/debug.h>
4  #include <tools/hash.h>
5  #include <tools/set.h>
6
7  #define CREATING_LLAMA_PARSER /* Suppress various definitions in parser.h */
8  #include "parser.h"           /* that conflict with llama-generated defs. */
9
10 /* This file is a llama input file that creates a parser for llama, like a snake
11 * eating its tail. The resulting yyparse.c file can be used in place of the
12 * recursive-descent parser in llpar.c. Note that, though this file is certainly
13 * easier to write than the recursive-descent version, the resulting code is
14 * about 1K larger. nows() should be called before firing up the parser.
15 * Most of the external subroutines called from this module are in acts.c.
16 * Exceptions are:
17 */
18
19 extern char      *yytext;          /* Generated by lex */
20 extern int       yylineno;        /* Generated by lex */
21 extern void      nows(), ws();    /* Declared in llama.lex. */
22
23 #define YYSTYPE char*           /* Value-stack type. */
24 %}
25
26 %token ACTION          /* {str} */
27 %token CODE_BLOCK       /* %{ ... %} */
28 %token COLON            /*
29 %token END_OPT          /* J J* */
30 %token FIELD             <name>
31 %token LEFT              %left

```



Listing 4.24. continued...

```

32 %token NAME          /*      name      */
33 %token NONASSOC      /*      %nonassoc   */
34 %token OR            /*      |      */
35 %token OTHER          /*      anything else */
36 %token PREC          /*      %prec     */
37 %token RIGHT          /*      %right    */
38 %token SEMI          /*      ;      */
39 %token SEPARATOR      /*      %%      */
40 %token START          /*      %start    */
41 %token START_OPT      /*      [      */
42 %token SYNCH          /*      %synch    */
43 %token TERM_SPEC      /*      %term or %token */
44 %token TYPE           /*      %type     */
45 %token PERCENT_UNION /*      %union    */
46 %token WHITESPACE     /*      0 <= c <= ' ' */
47
48 %synch SEMI OR
49 %%
50 spec      : defs SEPARATOR { first_sym(); } rules end
51 ;
52 end       : {ws();} SEPARATOR
53 | /* empty */
54 ;
55 defs      : SYNCH snames defs
56 | PERCENT_UNION ACTION { union_def( yytext ); } defs
57 | TYPE                  fnames { new_field(""); } defs
58 | TERM_SPEC  { new_lev ( 0 ); } tnames { new_field(""); } defs
59 | LEFT       { new_lev ('l'); } pnames { new_field(""); } defs
60 | RIGHT      { new_lev ('r'); } pnames { new_field(""); } defs
61 | NONASSOC   { new_lev ('n'); } pnames { new_field(""); } defs
62 | CODE_BLOCK  /* the block is copied out by yylex */  defs
63 | START
64 {
65     lerror( NONFATAL,
66             "%start not supported by occs. The first\n"
67             "\t\tt\tproduction is the start production\n");
68 }
69 opt_names defs
70
71 | /* empty */
72 ;
73 fnames    : NAME  { new_nonterm (yytext,0); } fnames
74 | FIELD { new_field  (yytext); } fnames
75 | /* empty */
76 ;
77 tnames    : NAME  { make_term(yytext); } tnames
78 | FIELD { new_field(yytext); } tnames
79 | /* empty */
80 ;
81 pnames    : NAME  { prec_list(yytext); } pnames
82 | FIELD { new_field(yytext); } pnames
83 | /* empty */
84 ;
85 opt_names : NAME opt_names
86 | /* empty */
87 ;
88 snames    : NAME { add_synch(yytext); } snames | /* empty */ ;
89 rules     : rule rules
90 | /* empty */
91 ;

```

Listing 4.24. continued...

```

92 rule           : NAME { new_nonterm(yytext,1); } COLON right_sides
93             | FIELD { new_nonterm(yytext,1); } COLON right_sides
94 ;
95 right_sides   : { new_rhs(); } rhs end_rhs
96 ;
97 end_rhs        : OR right_sides
98             | SEMI
99 ;
100 rhs           : NAME      { add_to_rhs(yytext, 0)           ); } rhs
101            | FIELD     { add_to_rhs(yytext, 0)           ); } rhs
102            | ACTION    { add_to_rhs(yytext, start_action()); } rhs
103            | PREC NAME { prec      (yytext)           ); } rhs
104            | START_OPT { start_opt( yytext ); }
105                  rhs END_OPT { end_opt  ( yytext ); }
106                  rhs
107            | /* empty */
108 ;
109 %%
110 yy_init_llama( tovs )
111 yyvstype *tovs;
112 {
113     tovs->left = tovs->right = "";
114 }
115
116 char *yypstk(tovs, tods)
117 yyvstype *tovs;
118 char *tods;
119 {
120     static char buf[128];
121
122     if( *tovs->left || *tovs->right )
123     {
124         sprintf(buf,"%s,%s", tovs->left, tovs->right);
125         return buf;
126     }
127     else
128         return "";
129 }
```

Listing 4.25. llpar.c— Recursive-Descent Parser for LLama

```

1 #include <stdio.h>
2 #include <stdarg.h>
3 #include <tools/debug.h>
4 #include <tools/set.h>
5 #include <tools/hash.h>
6 #include <tools/compiler.h>
7 #include <tools/l.h>
8
9 #include "llout.h"
10 #include "parser.h"
11
12 /* LLPAR.C:      A recursive-descent parser for a very stripped down llama.
13  *                  There's a llama input specification for a table-driven parser
14  *                  in llama.lma.
15  */
16
```

Listing 4.25. continued...

```

17 void advance      P(( void ));          /* local */
18 void lookfor      P(( int first,... )); 
19 void definitions  P(( void )) ;
20 void body         P(( void )) ;
21 void right_sides P(( void )) ;
22 void rhs          P(( void )) ;
23 int yyparse       P(( void )) ;        /* public */
24 extern int        yylineno;           /* Created by lex.*/ 
25 extern char       *yytext;
26 extern int        yylex(); 
27 /*-----*/
28
29 PUBLIC int yynerrs;          /* Total error count */ 
30 PRIVATE int Lookahead;        /* Lookahead token */ 
31
32 /*=====
33  * Low-level support routines for parser:
34  */
35
36 #define match(x) ((x) == Lookahead)
37
38 PRIVATE void advance()
39 {
40     if( Lookahead != _EOI_ )
41         while( (Lookahead = yylex()) == WHITESPACE )
42             ;
43 }
44
45 /*-----*/
46
47 PRIVATE void    lookfor( first, ... )
48 int      first;
49 {
50     /* Read input until the current symbol is in the argument list. For example,
51      * lookfor(OR, SEMI, 0) terminates when the current Lookahead symbol is an
52      * OR or SEMI. Searching starts with the next (not the current) symbol.
53      */
54
55     int *obj;
56
57     for( advance() ;; advance() )
58     {
59         for( obj = &first; *obj && !match(*obj) ; obj++ )
60             ;
61
62         if( *obj )
63             break;           /* Found a match */
64
65         else if( match(_EOI_) )
66             lerror(FATAL, "Unexpected end of file\n");
67     }
68 }
69
70 /*=====
71  * The Parser itself
72  */
73
74 PUBLIC int yyparse()          /* spec : definitions body stuff */ 
75 {
76     extern yylineno;

```

Listing 4.25. continued...

```

77     Lookahead = yylex();           /* Get first input symbol */
78     definitions();;
79     first_sym();
80     body();
81 }
82
83 /*-----*/
84
85 PRIVATE void    definitions()
86 {
87     /*
88      * definitions  : TERM_SPEC tnames  definitions
89      *                 | CODE_BLOCK      definitions
90      *                 | SYNCH snames     definitions
91      *                 | SEPARATOR
92      *                 | _EOI_
93      * snames       : NAME {add_synch} snames
94      * tnames       : NAME {make_term} tnames
95
96      * Note that LeX copies the CODE_BLOCK contents to the output file
97      * automatically on reading it.
98     */
99
100    while( !match(SEPARATOR) && !match(_EOI_) )          /* 4 */
101    {
102        if( Lookahead == SYNCH )                         /* 3 */
103        {
104            for( advance(); match(NAME); advance() )      /* 5 */
105                add_synch( yytext );
106        }
107        else if( Lookahead == TERM_SPEC )               /* 1 */
108        {
109            for( advance(); match(NAME); advance() )      /* 6 */
110                make_term( yytext );
111        }
112        else if( Lookahead != CODE_BLOCK )              /* 2 */
113        {
114            lerror(NONFATAL,"Ignoring illegal <%s> in definitions\n", yytext );
115            advance();
116        }
117    }
118
119    advance();                                         /* advance past the %% */
120 }
121
122 /*-----*/
123
124 PRIVATE void    body()
125 {
126     /*
127      * body      : rule body
128      *                 | rule SEPARATOR
129      *                 | rule _EOI_
130      * rule      : NAME {new_nonterm} COLON right_sides
131      *                 : <epsilon>
132     */
133
134    while( !match(SEPARATOR) && !match(_EOI_) )          /* 1 */
135    {

```

implemented at:

1	
2	
3	
4	
5	
6	



Listing 4.25. continued...

```

136         if( match(NAME) )                                /* 2 */
137         {
138             new_nonterm( yytext, 1 );
139             advance();
140         }
141     else                                              /* 3 */
142     {
143         lerror(NONFATAL, "Illegal <%s>, nonterminal expected.\n", yytext);
144         lookfor( SEMI, SEPARATOR, 0 );
145         if( match(SEMI) )
146             advance();
147         continue;
148     }
149
150     if( match(COLON) )
151         advance();
152     else
153         lerror(NONFATAL, "Inserted missing ':'\n");
154
155     right_sides();
156 }
157 ws();                                         /* Enable white space (see parser.lex) */
158 if( match(SEPARATOR) )
159     yylex();                                     /* Advance past %% */
160 }
161
162 /*-----*/
163
164 PRIVATE void      right_sides()
165 {
166     /* right_sides      : {new_rhs} rhs OR right_sides
167      *                  | {new_rhs} rhs SEMI
168      */
169
170     new_rhs();
171     rhs();
172
173     while( match(OR) )                           /* 1 */
174     {
175         advance();
176         new_rhs();
177         rhs();
178     }
179
180     if( match(SEMI) )                          /* 2 */
181         advance();
182     else
183         lerror(NONFATAL, "Inserted missing semicolon\n");
184 }
185
186 /*-----*/
187
188 PRIVATE void      rhs()
189 {
190     /* rhs      : NAME    {add_to_rhs} rhs   */
191     /*          | ACTION  {add_to_rhs} rhs   */
192
193     while( match(NAME) || match(ACTION) )
194     {
195         add_to_rhs( yytext, match(ACTION) ? start_action() : 0 );

```



Listing 4.25. continued...

```

196         advance();
197     }
198
199     if( !match(OR) && !match(SEMI) )
200     {
201         lerror(NONFATAL, "illegal <%s>, ignoring rest of production\n", yytext);
202         lookfor( SEMI, SEPARATOR, OR, 0 );
203     }
204 }
```

The action subroutines, in Listings 4.26, 4.27, and 4.28. work as follows:

LLama actions:

`add_synch(char *yytext)`

Listing 4.28, lines 628 and 690.

Add the symbol named by `yytext` to the synchronization set used for error recovery.

`add_to_rhs(char *yytext, int isact)`

Listing 4.27, line 488.

Add the symbol named by `yytext` to the current right-hand side. The `isact` argument should be true if the symbol is an action.

`, end_opt(char *yytext)`

Listing 4.28, line 596.

Mark the end of an optional production. This routine processes `occ's`' [...] , [...] * operators by adding extra elements to the production.

`first_sym(void)`

Listing 4.27, line 380.

This routine finds the goal symbol, which is the left-hand side of the first production that follows the first % directive. The problem here is that, in `occ's`, a nonterminal can be declared explicitly in a previous %type directive, so you can't assume that the first nonterminal that's declared is the goal symbol.

`lerror(int error_type, char *format,...)`

In main.c.

This routine prints an error message. The first argument should be FATAL, NON-FATAL, or WARNING (these are defined at the top of *parser.h*). If the error is FATAL, all files are closed and LLama is terminated, otherwise parsing continues. The input file name and line number is automatically appended to the front of the message. This routine works like `fprintf()` in other respects.

`make_term(char *yytext)`

Listing 4.27, line 335.

Make an entry for a terminal symbol.

`new_field(char *yytext)`

Listing 4.28, lines 675 and 793.

Change the name of the %union field to which subsequent input symbols are attached.

`new_lev(int associativity)`

Listing 4.28, lines 648 and 698.

Increment the precedence level by one notch (all subsequent arguments to a %left, %right, or %nonassoc use this new level). Also change the current associativity as specified. This argument should be 'l' for left , 'r' for right, 'n' for none, or 0 for unspecified.

Listing 4.27, line 391.

```
new_nonterm(char *yytext, int is_lhs)
```

Create a symbol-table entry for a nonterminal if one doesn't already exist. yytext is the symbol name. is_lhs tells the subroutine whether the nonterminal was created implicitly (whether it was used before appearing on the left-hand side of a production). is_lhs zero in this case. Return a pointer to the SYMBOL or NULL if an attempt is made to use a terminal symbol as a left-hand side.

Listing 4.27, line 460.

```
new_rhs(void)
```

Stop adding symbols to the current right-hand side, and start up a new (and empty) right-hand side.

In main.c.

```
output(char *format, ...)
```

Works like `fprintf(Output, ...);` sends the string to the current output stream, typically standard output but it can be changed from the command line.

Listing 4.28, lines 659 and 735.

```
prec(char *yytext)
```

Process the argument to a %prec directive. (Change the precedence level of the current right-hand side from the default level to the one indicated in yytext.)

Listing 4.28, lines 671 and 710.

```
prec_list(char *yytext)
```

Process an argument to %left, %right, or %nonassoc.

Listing 4.28, line 581.

```
start_opt(char *yytext)
```

Mark the start of an optional or repeating portion of a right-hand side [see end_opt()].

Listing 4.28, lines 665 and 762.

```
union_def(char *yytext)
```

Create a typedef for the occs value stack, using the %union definition in yytext.

Listing 4.26. *acts.c*—LLama Action Subroutines

```

1  #include <stdio.h>
2  #include <malloc.h>
3  #include <ctype.h>
4  #include <string.h>
5  #include <tools/debug.h>
6  #include <tools/set.h>
7  #include <tools/hash.h>
8  #include <tools/compiler.h>
9  #include <tools/l.h>
10
11 #include <tools/stack.h>           /* stack-manipulation macros */
12 #undef stack_cls                /* Make all stacks static */
13 #define stack_cls static
14
15 #include "parser.h"
16 #include "llout.h"
17
18 /* ACTS.C      Action routines used by both llama and occs. These build
19 *                  up the symbol table from the input specification.
20 */
21

```



Listing 4.26. continued...

```

22 void find_problems P(( SYMBOL *sym, void *junk ))/* local */;
23 int c_identifier P(( char *name ));;
24 void print_tok P(( FILE *stream, char *format, int arg ));/* public */;
25 void pterm P(( SYMBOL *sym, FILE *stream ));;
26 void pact P(( SYMBOL *sym, FILE *stream ));;
27 void pnonterm P(( SYMBOL *sym, FILE *stream ));;
28 void print_symbols P(( FILE *stream ));;
29 int problems P(( void ));;
30 void init_acts P(( void ));;
31 SYMBOL *make_term P(( char *name ));;
32 void first_sym P(( void ));;
33 SYMBOL *new_nonterm P(( char *name, int is_lhs ));;
34 void new_rhs P(( void ));;
35 void add_to_rhs P(( char *object, int is_an_action ));;
36 void start_opt P(( char *lex ));;
37 void end_opt P(( char *lex ));;
38 void add_synch P(( char *yytext ));;
39 void new_lev P(( int how ));;
40 void prec_list P(( char *name ));;
41 void prec P(( char *name ));;
42 void union_def P(( char *action ));;
43 int fields_active P(( void ));;
44 void new_field P(( char *field_name ));;
45
46 /*-----*/
47
48 extern int yylineno; /* Input line number--created by Lex. */
49 PRIVATE int Associativity; /* Current associativity direction. */
50 PRIVATE int Prec_lev = 0; /* Precedence level. Incremented */
51 /* after finding %left, etc., */
52 /* but before the names are done. */
53 PRIVATE char Field_name[NAME_MAX]; /* Field name specified in <name>. */
54 PRIVATE int Fields_active = 0; /* Fields are used in the input. */
55 /* (If they're not, then automatic */
56 /* field-name generation, as per */
57 /* %union, is not activated.) */
58 PRIVATE int Goal_symbol_is_next = 0; /* If true, the next nonterminal is */
59 /* the goal symbol. */
60
61 /*-----*/
62 /* The following stuff (that's a technical term) is used for processing nested */
63 /* optional (or repeating) productions defined with the occs [] and []* */
64 /* operators. A stack is kept and, every time you go down another layer of */
65 /* nesting, the current nonterminal is stacked and a new nonterminal is */
66 /* allocated. The previous state is restored when you're done with a level. */
67 */
68
69 #define SSIZE 8 /* Max. optional-production nesting level */
70 typedef struct _cur_sym_
71 {
72     char lhs_name[NAME_MAX]; /* Name associated with left-hand side. */
73     SYMBOL *lhs; /* Pointer to symbol-table entry for */
74                                         /* the current left-hand side. */
75     PRODUCTION *rhs; /* Pointer to current production. */
76
77 } CUR_SYM;
78
79 CUR_SYM Stack[ SSIZE ], /* Stack and */
80             *Sp = Stack + (SSIZE-1); /* stack pointer. It's inconvenient to use */

```



Listing 4.26. continued...

```

81                                     /* stack.h because stack is of structures. */
82
83 /*=====
84 * Support routines for actions
85 */
86
87 PUBLIC void print_tok( stream, format, arg )
88 FILE    *stream;
89 char    *format;           /* not used here but supplied by pset() */
90 int     arg;
91 {
92     /* Print one nonterminal symbol to the specified stream. */
93
94     if      ( arg == -1      ) fprintf(stream, "null " );
95     else if( arg == -2      ) fprintf(stream, "empty " );
96     else if( arg == _EOI_   ) fprintf(stream, "$ " );
97     else if( arg == EPSILON ) fprintf(stream, "<epsilon> " );
98     else
99         fprintf(stream, "%s ", Terms[arg]->name );
100 }
101 /*
102 * The following three routines print the symbol table. Pterm(), pact(),
103 * pnonterm() are all called indirectly through ptab(), called in
104 * print_symbols(). They print the terminal, action, and nonterminal symbols
105 * from the symbol table, respectively.
106 */
107
108 PUBLIC void pterm( sym, stream )
109 SYMBOL *sym;
110 FILE   *stream;
111 {
112     int i;
113
114     if( !ISTERM(sym) )
115         return;
116
117     LL( fprintf( stream, "%-16.16s %3d\n", sym->name, sym->val ) );
118     OX( fprintf( stream, "%-16.16s %3d %2d %c <%s>\n", \
119          sym->name, \
120          sym->val, \
121          Precedence[sym->val].level , \
122          (i = Precedence[sym->val].assoc) ? i : '-', \
123          sym->field ); )
124 }
125
126 /*
127
128 PUBLIC void pact( sym, stream )
129 SYMBOL *sym;
130 FILE   *stream;
131 {
132     if( !ISACT(sym) )
133         return;
134
135     fprintf( stream, "%-5s %3d,",     sym->name, sym->val );
136     fprintf( stream, " line %-3d: ", sym->lineno );
137     fputstr( sym->string, 55, stream );
138     fprintf( stream, "\n" );
139 }
140

```



Listing 4.26. continued...

```

141  /*-----*/
142
143 PUBLIC char *production_str( prod )
144 PRODUCTION *prod;
145 {
146     /* return a string representing the production */
147
148     int      i, nchars, avail;
149     static char buf[80];
150     char     *p;
151
152     nchars = sprintf(buf,"%s ->", prod->lhs->name );
153     p      = buf + nchars;
154     avail  = sizeof(buf) - nchars - 1;
155
156     if( !prod->rhs_len )
157         sprintf(p, " (epsilon)" );
158     else
159         for( i = 0; i < prod->rhs_len && avail > 0 ; ++i )
160         {
161             nchars = sprintf(p, " %0.*s", avail-2, prod->rhs[i]->name );
162             avail -= nchars;
163             p    += nchars;
164         }
165
166     return buf;
167 }
168
169 /*-----*/
170
171 PUBLIC void pnonterm( sym, stream )
172 SYMBOL *sym;
173 FILE   *stream;
174 {
175     int      i;
176     char    *s;
177     PRODUCTION *p;
178     int      chars_printed;
179
180     stack_dcl( pstack, PRODUCTION *, MAXPROD );
181
182     if( !ISNONTERM(sym) )
183         return;
184
185     fprintf( stream, "%s (%3d)  %s", sym->name, sym->val,
186             sym == Goal_symbol ? "(goal symbol)" : "" );
187
188     OX( fprintf( stream, " <%s>\n", sym->field );           )
189     LL( fprintf( stream, "\n" );                                )
190
191     if( Symbols > 1 )
192     {
193         /* Print first and follow sets only if you want really verbose output.*/
194
195         fprintf( stream, " FIRST : " );
196         pset( sym->first, print_tok, stream );
197
198         LL( fprintf(stream, "\n FOLLOW: " );
199         LL( pset( sym->follow, print_tok, stream );      )
200

```



Listing 4.26. continued...

```

201         fprintf(stream, "\n");
202     }
203
204     /* Productions are put into the SYMBOL in reverse order because it's easier
205      * to tack them on to the beginning of the linked list. It's better to print
206      * them in forward order, however, to make the symbol table more readable.
207      * Solve this problem by stacking all the productions and then popping
208      * elements to print them. Since the pstack has MAXPROD elements, it's not
209      * necessary to test for stack overflow on a push.
210     */
211
212     for( p = sym->productions ; p ; p = p->next )
213         push( pstack, p );
214
215     while( !stack_empty( pstack ) )
216     {
217         p = pop(pstack);
218
219         chars_printed = fprintf(stream, "    %3d: %s",
220                               p->num, production_str( p ));
221
222         LL( for( ; chars_printed <= 45; ++chars_printed ) )
223         LL(     putc( '.', stream ) );
224         LL(     fprintf(stream, "SELECT: " );
225         LL(     pset( p->select, print_tok, stream ) );
226
227         OX( if( p->prec ) )
228         OX( {
229             OX(     for( ; chars_printed <= 60; ++chars_printed ) )
230             OX(                 putc( '.', stream ) );
231             OX(                 if( p->prec ) )
232                 OX(                     fprintf(stream, "PREC %d", p->prec ) );
233             OX(             )
234
235             putc('\n', stream);
236         }
237
238         fprintf(stream, "\n");
239     }
240
241     /*-----*/
242
243 PUBLIC void print_symbols( stream )
244 FILE    *stream;
245 {
246     /* Print out the symbol table. Nonterminal symbols come first for the sake
247      * of the 's' option in yydebug(); symbols other than production numbers can
248      * be entered symbolically. ptab returns 0 if it can't print the symbols
249      * sorted (because there's no memory. If this is the case, try again but
250      * print the table unsorted).
251     */
252
253     putc( '\n', stream );
254
255     fprintf( stream, "----- Symbol table -----\\n" );
256     fprintf( stream, "\\nNONTERMINAL SYMBOLS:\\n\\n" );
257     if( ptab( Symtab, pnonterm, stream, 1 ) == 0 )
258         ptab( Symtab, pnonterm, stream, 0 );
259

```



Listing 4.26. continued...

```

260     fprintf( stream, "\nTERMINAL SYMBOLS:\n\n");
261     OX( fprintf( stream, "name           value  prec  assoc   field\n"); )
262     LL( fprintf( stream, "name           value\n"); )
263
264     if( ptab( Symtab, pterm, stream, 1 ) == 0 )
265         ptab( Symtab, pterm, stream, 0 );
266
267     LL( fprintf( stream, "\nACTION SYMBOLS:\n\n");
268     LL( if( !ptab( Symtab, pact, stream, 1 ) )
269         LL( ptab( Symtab, pact, stream, 0 );
270     LL( fprintf( stream, "-----\n" ); )
271 }
272
273 /*-----
274 * Problems() and find_problems work together to find unused symbols and
275 * symbols that are used but not defined.
276 */
277
278 PRIVATE void find_problems( sym, junk )
279 SYMBOL *sym;
280 void *junk; /* not used */
281 {
282     if( !sym->used && sym!=Goal_symbol )
283         error( WARNING, "<%s> not used (defined on line %d)\n",
284                     sym->name, sym->set );
285     if( !sym->set && !ISACT(sym) )
286         error( NONFATAL, "<%s> not defined (used on line %d)\n",
287                     sym->name, sym->used );
288 }
289
290 PUBLIC int problems()
291 {
292     /* Find, and print an error message, for all symbols that are used but not
293      * defined, and for all symbols that are defined but not used. Return the
294      * number of errors after checking.
295 */
296
297     ptab( Symtab, find_problems, NULL, 0 );
298     return yynerrs;
299 }
300
301 /*-----*/
302
303 PRIVATE int hash_funct( p )
304 SYMBOL *p;
305 {
306     if( !*p->name )
307         lerror( FATAL, "Illegal empty symbol name\n" );
308
309     return hash_add( p->name );
310 }
311
312 PUBLIC void init_acts()
313 {
314     /* Various initializations that can't be done at compile time. Call this
315      * routine before starting up the parser. The hash-table size (157) is
316      * an arbitrary prime number, roughly the number symbols expected in the
317      * table.
318 */
319 }
```



Listing 4.26. continued...

```

320     extern int      strcmp();
321     static SYMBOL bogus_symbol;
322
323     strcpy( bogus_symbol.name, "End of Input" );
324
325     Terms[0]      = &bogus_symbol;
326     Symtab       = maketab( 157, hash_funct, strcmp );
327
328     strncpy( Sp->lhs_name, "garbage", NAME_MAX );
329
330     LL( Synch    = newset(); )
331 }
332
333 /*-----*/
334
335 PUBLIC SYMBOL *make_term( name )           /* Make a terminal symbol */
336 char      *name;
337 {
338     SYMBOL      *p;
339
340     if( !c_identifier(name) )
341         lerror(NONFATAL, "Token names must be legitimate C identifiers\n");
342
343     else if( p = findsym(Symtab, name) )
344         lerror(WARNING, "Terminal symbol <%s> already declared\n", name );
345     else
346     {
347         if( Cur_term >= MAXTERM )
348             lerror(FATAL, "Too many terminal symbols (%d max.).\n", MAXTERM );
349
350         p = (SYMBOL *) newsym( sizeof(SYMBOL) );
351         strncpy ( p->name, name,          NAME_MAX );
352         strncpy ( p->field, Field_name, NAME_MAX );
353         addsym ( Symtab, p );
354
355         p->val = ++Cur_term ;
356         p->set = yylineno;
357
358         Terms[Cur_term] = p;
359     }
360     return p;
361 }
362
363 /*-----*/
364
365 PRIVATE c_identifier( name )           /* Return true only if name is */
366 char      *name;                      /* a legitimate C identifier. */
367 {
368     if( isdigit( *name ) )
369         return 0;
370
371     for( *name ; ++name )
372         if( !( isalnum(*name) || *name == '_' ))
373             return 0;
374
375     return 1;
376 }
377

```



Listing 4.26. continued...

```

378  /*-----*/
379
380 PUBLIC void      first_sym()
381 {
382     /* This routine is called just before the first rule following the
383      * %. It's used to point out the goal symbol;
384      */
385
386     Goal_symbol_is_next = 1;
387 }
388
389 /*-----*/
390
391 PUBLIC SYMBOL  *new_nonterm( name, is_lhs )
392 char   *name;
393 int    is_lhs;
394 {
395     /* Create, and initialize, a new nonterminal. is_lhs is used to
396      * differentiate between implicit and explicit declarations. It's 0 if the
397      * nonterminal is added because it was found on a right-hand side. It's 1 if
398      * the nonterminal is on a left-hand side.
399      */
400
401     /* Return a pointer to the new symbol or NULL if an attempt is made to use a
402      * terminal symbol on a left-hand side.
403      */
404
405     SYMBOL      *p;
406
407     if( p = (SYMBOL *) findsym( Symtab, name ) )
408     {
409         if( !ISNONTERM( p ) )
410         {
411             lerror(NONFATAL, "Symbol on left-hand side must be nonterminal\n");
412             p = NULL;
413         }
414     else if( Cur_nonterm >= MAXNONTERM )
415     {
416         lerror(FATAL, "Too many nonterminal symbols (%d max.).\n", MAXTERM );
417     }
418     else
419         /* Add new nonterminal to symbol table */
420
421         p = (SYMBOL *) newsym( sizeof(SYMBOL) );
422         strncpy ( p->name, name, NAME_MAX );
423         strncpy ( p->field, Field_name, NAME_MAX );
424
425         p->val = ++Cur_nonterm ;
426         Terms[Cur_nonterm] = p;
427
428         addsym ( Symtab, p );
429     }
430
431     if( p )
432         /* (re)initialize new nonterminal */
433     {
434         if( Goal_symbol_is_next )
435         {
436             Goal_symbol = p;
437             Goal_symbol_is_next = 0;
438         }

```



Listing 4.26. continued...

```

438         if( !p->first )
439             p->first    = newset();
440
441     LL(   if( !p->follow )           )
442     LL(       p->follow = newset();      )
443
444         p->lineno = yylineno ;
445
446         if( is_lhs )
447         {
448             strncpy( Sp->lhs_name, name, NAME_MAX );
449             Sp->lhs      = p ;
450             Sp->rhs      = NULL;
451             Sp->lhs->set = yylineno;
452         }
453     }
454
455     return p;
456 }
457
458 /*-----*/
459
460 PUBLIC void new_rhs()
461 {
462     /* Get a new PRODUCTION and link it to the head of the production chain.
463      * of the current nonterminal. Note that the start production MUST be
464      * production 0. As a consequence, the first rhs associated with the first
465      * nonterminal MUST be the start production. Num_productions is initialized
466      * to 0 when it's declared.
467     */
468
469     PRODUCTION *p;
470
471     if( !(p = (PRODUCTION *) calloc(1, sizeof(PRODUCTION))) )
472         lerror(FATAL, "Out of memory\n");
473
474     p->next          = Sp->lhs->productions;
475     Sp->lhs->productions = p;
476
477     LL( p->select = newset(); )
478
479     if( (p->num = Num_productions++) >= MAXPROD )
480         lerror(FATAL, "Too many productions (%d max.)\n", MAXPROD );
481
482     p->lhs  = Sp->lhs;
483     Sp->rhs = p;
484 }
485
486 /*-----*/
487
488 PUBLIC void add_to_rhs( object, is_an_action )
489 char      *object;
490 int       is_an_action; /* 0 if not an action, line number otherwise */
491 {
492     SYMBOL      *p;
493     PRODUCTION *prod;
494     int i;
495     char       buf[32];
496

```



Listing 4.26. continued...

```

497     /* Add a new element to the RHS currently at top of stack. First deal with
498      * forward references. If the item isn't in the table, add it. Note that,
499      * since terminal symbols must be declared with a %term directive, forward
500      * references always refer to nonterminals or action items. When we exit the
501      * if statement, p points at the symbol table entry for the current object.
502      */
503
504     if( !(p = (SYMBOL *) findsym( Symtab, object)) ) /* not in tab yet */
505     {
506         if( !is_an_action )
507         {
508             if( !(p = new_nonterm( object, 0 )) )
509             {
510                 /* Won't get here unless p is a terminal symbol */
511
512                 lerror(FATAL, "(internal) Unexpected terminal symbol\n");
513                 return;
514             }
515         }
516     else
517     {
518         /* Add an action. All actions are named "{DDD}" where DDD is the
519          * action number. The curly brace in the name guarantees that this
520          * name won't conflict with a normal name. I am assuming that calloc
521          * is used to allocate memory for the new node (ie. that it's
522          * initialized to zeros).
523         */
524
525         sprintf(buf, "%d", ++Cur_act - MINACT );
526
527         p = (SYMBOL *) newsym( sizeof(SYMBOL) );
528         strncpy ( p->name, buf, NAME_MAX );
529         addsym ( Symtab, p );
530
531         p->val      = Cur_act      ;
532         p->lineno = is_an_action ;
533
534         if( !(p->string = strsave(object)) )
535             lerror(FATAL, "Insufficient memory to save action\n");
536     }
537 }
538
539 p->used = yylineno;
540
541 if( (i = Sp->rhs->rhs_len++) >= MAXRHS )
542     lerror(NONFATAL, "Right-hand side too long (%d max)\n", MAXRHS );
543 else
544 {
545     LL( if( i == 0 && p == Sp->lhs ) )
546     LL(     lerror(NONFATAL, "Illegal left recursion in production.\n"); )
547     OX( if( ISTERM( p ) ) )
548     OX(     Sp->rhs->prec = Precedence[ p->val ].level ; )
549
550     Sp->rhs->rhs[ i ] = p;
551     Sp->rhs->rhs[ i + 1 ] = NULL;           /* NULL terminate the array. */
552
553     if( !ISACT(p) )
554         ++( Sp->rhs->non_acts );
555 }
556 }
```



Listing 4.26. continued...

```

557
558 /* -----
559 * The next two subroutines handle repeating or optional subexpressions. The
560 * following mappings are done, depending on the operator:
561 *
562 * S : A [B] C ;           S    -> A 001 C
563 *                   001 -> B | epsilon
564 *
565 * S : A [B]* C ;        S    -> A 001 C          (occs)
566 *                   001 -> 001 B | epsilon
567 *
568 * S : A [B]* C ;        S    -> A 001 C          (llama)
569 *                   001 -> B 001 | epsilon
570 *
571 * In all situations, the right hand side that we've collected so far is
572 * pushed and a new right-hand side is started for the subexpression. Note that
573 * the first character of the created rhs name (001 in the previous examples)
574 * is a space, which is illegal in a user-supplied production name so we don't
575 * have to worry about conflicts. Subsequent symbols are added to this new
576 * right-hand side. When the ), ], or *) is found, we finish the new right-hand
577 * side, pop the stack and add the name of the new right-hand side to the
578 * previously collected left-hand side.
579 */
580
581 PUBLIC void start_opt( lex )           /* Start an optional subexpression */
582 char   *lex;
583 {
584     char   name[32], *tname;
585     static int num = 0;
586
587     --Sp;                                /* Push current stack element */
588     sprintf( name, " %06d", num++ );
589     new_nonterm( name, 1 );
590     new_rhs();
591     new_rhs();                            /* Make name for new production */
592 }                                         /* Create a nonterminal for it. */
593
594 /*----- */
595
596 PUBLIC void end_opt( lex )             /* end optional subexpression */
597 char   *lex;
598 {
599     char   *name = Sp->lhs_name ;
600     SYMBOL *p;
601     int    i;
602
603     if( lex[1] == '*' )                  /* Process a [...] */
604     {
605         add_to_rhs( name, 0 );          /* Add right-recursive reference. */
606
607 #ifdef OCCS
608         i = Sp->rhs->rhs_len - 1;      /* If occs, must be left recursive.*/
609         p = Sp->rhs->rhs[ i ];
610         memmove( &(Sp->rhs->rhs)[1], &(Sp->rhs->rhs)[0],
611                   i * sizeof( (Sp->rhs->rhs)[1] ) );
612         Sp->rhs->rhs[ 0 ] = p ;
613 #endif
614     }
615

```



Listing 4.26. continued...

```

616     ++Sp;                                /* discard top-of-stack element */
617     add_to_rhs( name, 0 );
618 }
619
620 /*=====
621 * The following routines have alternate versions, one set for llama and another
622 * for occs. The routines corresponding to features that aren't supported in one
623 * or the other of these programs print error messages.
624 */
625
626 #ifdef LLAMA
627
628 PUBLIC void add_synch( name )
629 char    *name;
630 {
631     /* Add "name" to the set of synchronization tokens
632      */
633
634     SYMBOL      *p;
635
636     if( !(p = (SYMBOL *) findsym( Symtab, name )) )
637         lerror(NONFATAL,"%synch: undeclared symbol <%s>.\n", name );
638
639     else if( !ISTERM(p) )
640         lerror(NONFATAL,"%synch: <%s> not a terminal symbol\n", name );
641
642     else
643         ADD( Synch, p->val );
644 }
645
646 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
647
648 PUBLIC void new_lev( how )
649 {
650     switch( how )
651     {
652         case 0 : /* initialization: ignore it */                                break;
653         case 'l': lerror (NONFATAL, "%left not recognized by LLAMA\n" ); break;
654         case 'r': lerror (NONFATAL, "%right not recognized by LLAMA\n" ); break;
655         default : lerror (NONFATAL, "%nonassoc not recognized by LLAMA\n"); break;
656     }
657 }
658
659 PUBLIC void prec( name )
660 char    *name;
661 {
662     lerror( NONFATAL, "%prec not recognized by LLAMA\n" );
663 }
664
665 PUBLIC void union_def( action )
666 char    *action;
667 {
668     lerror(NONFATAL,"%union not recognized by LLAMA\n");
669 }
670
671 PUBLIC void prec_list( name ) char *name;
672 {
673 }
674

```



Listing 4.26. continued...

```

675 PUBLIC void new_field( field_name )
676     char *field_name;
677 {
678     if( *field_name )
679         lerror(NONFATAL, "<name> not supported by LLAMA\n");
680 }
681
682 PUBLIC make_nonterm(name)
683     char *name;
684 {
685     lerror(NONFATAL, "%type not supported by LLAMA\n");
686 }
687
688 #else /*=====
689
690 PUBLIC void add_synch(yytext)
691     char *yytext;
692 {
693     lerror(NONFATAL, "%synch not supported by OCCS\n");
694 }
695
696 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
697
698 PUBLIC void new_lev( how )
699 {
700     /* Increment the current precedence level and modify "Associativity"
701      * to remember if we're going left, right, or neither.
702      */
703
704     if( Associativity = how ) /* 'l', 'r', 'n', (0 if unspecified) */
705         ++Prec_lev;
706 }
707
708 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
709
710 PUBLIC void prec_list( name )
711     char *name;
712 {
713     /* Add current name (in yytext) to the precision list. "Associativity" is
714      * set to 'l', 'r', or 'n', depending on whether we're doing a %left,
715      * %right, or %nonassoc. Also make a nonterminal if it doesn't exist
716      * already.
717      */
718
719     SYMBOL *sym;
720
721     if( !(sym = findsym(Syntab, name)) )
722         sym = make_term( name );
723
724     if( !ISTERM(sym) )
725         lerror(NONFATAL, "%left or %right, %s must be a token\n", name );
726     else
727     {
728         Precedence[ sym->val ].level = Prec_lev ;
729         Precedence[ sym->val ].assoc = Associativity ;
730     }
731 }
732

```



Listing 4.26. continued...

```

733 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
734
735 PUBLIC void prec( name )
736 char *name;
737 {
738     /* Change the precedence level for the current right-hand side, using
739      * (1) an explicit number if one is specified, or (2) an element from the
740      * Precedence[] table otherwise.
741      */
742
743     SYMBOL      *sym;
744
745     if( isdigit(*name) )
746         Sp->rhs->prec = atoi(name);                                /* (1) */
747     else
748     {
749         if( !(sym = findsym(Syntab,name)) )
750             lerror(NONFATAL, "%s (used in %%prec) undefined\n" );
751
752         else if( !ISTERM(sym) )
753             lerror(NONFATAL, "%s (used in %%prec) must be terminal symbol\n" );
754
755         else
756             Sp->rhs->prec = Precedence[ sym->val ].level;          /* (2) */
757     }
758 }
759
760 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
761
762 PUBLIC void union_def( action )
763 char *action;
764 {
765     /* create a YYSTYPE definition for the union, using the fields specified
766      * in the %union directive, and also appending a default integer-sized
767      * field for those situation where no field is attached to the current
768      * symbol.
769      */
770
771     while( *action && *action != '{' ) /* Get rid of everything up to the */
772         ++action;                      /* open brace */                  */
773     if( *action )                   /* and the brace itself */
774         ++action;
775
776     output( "typedef union\n" );
777     output( "{\n" );
778     output( "    int    %s; /* Default field, used when no %%type found */",
779                         DEF_FIELD );
780     output( "%s\n", action );
781     output( "yystype;\n\n" );
782     output( "#define YYSTYPE yystype\n" );
783     Fields_active = 1;
784 }
785
786 PUBLIC fields_active()
787 {
788     return Fields_active;           /* previous %union was specified */
789 }
790

```



Listing 4.26. continued...

```

791 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
792
793 PUBLIC void new_field( field_name )
794 char *field_name;
795 {
796     /* Change the name of the current <field> */
797
798     char *p;
799
800     if( !*field_name )
801         *Field_name = '\0' ;
802     else
803     {
804         if( p = strchr(++field_name, '>') )
805             *p = '\0' ;
806
807         strncpy( Field_name, field_name, sizeof(Field_name) );
808     }
809 }
810 #endif

```

4.10.2 Creating The Tables

The next step is to create the parse tables from the grammar we just assembled, and the LL(1) select sets are needed to do this.

4.10.2.1 Computing FIRST, FOLLOW, and SELECT Sets. **LLama** starts by creating the FIRST sets, using the code in Listing 4.27, which implements the algorithm described earlier. The subroutine `first()`, on line 21, figures the FIRST sets for all the nonterminals in the symbol table. The `ptab()` call on line 30 traverses the table, calling `first_closure()` (on line 37) for every table element (`ptab()` is one of the hash functions described in Appendix A). Multiple passes are made through the table, until nothing is added to any symbol's FIRST set.

The `first_rhs()` subroutine on line 97 computes the first set for an entire right-hand side, represented as an array of SYMBOL pointers. (They're stored this way in a PRODUCTION structure). `first_rhs()` won't work if called before `first()`.

The subroutines in Listing 4.28 compute FOLLOW sets for all nonterminals in **LLama**'s symbol table, again using the procedure described earlier in this Chapter. Finally, **LLama** uses the FIRST and FOLLOW sets to create the LL(1) selection sets in Listing 4.29.

4.10.3 The Rest of LLama

The remainder of **LLama** is Listings 4.30. to 4.37. All that they do is print out the tables, and so forth. They are adequately commented and need no further description here.

Listing 4.27. *first.c*—Find FIRST Sets

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/set.h>
4 #include <tools/hash.h>
5 #include <tools/compiler.h>
6 #include <tools/l.h>
7 #include "parser.h"
8
9 /* FIRST.C      Compute FIRST sets for all productions in a symbol table.
10  *-----*/
11 */
12
13 int first_closure P(( SYMBOL *lhs ));           /* local */
14 void first      P(( void ));                   /* public */
15 int first_rhs   P(( SET *dest, SYMBOL **rhs, int len ));
16
17 PRIVATE int Did_something;
18
19 /*-----*/
20
21 PUBLIC void first( )
22 {
23     /* Construct FIRST sets for all nonterminal symbols in the symbol table. */
24
25     D( printf("Finding FIRST sets\n"); )
26
27     do
28     {
29         Did_something = 0;
30         ptab( Symtab, first_closure, NULL, 0 );
31
32     } while( Did_something );
33 }
34
35 /*-----*/
36
37 PRIVATE first_closure( lhs )
38 SYMBOL *lhs;                                     /* Current left-hand side */
39 {
40     /* Called for every element in the FIRST sets. Adds elements to the first
41      * sets. The following rules are used:
42      *
43      * 1) given lhs->...Y... where Y is a terminal symbol preceded by any number
44      *      (including 0) of nullable nonterminal symbols or actions, add Y to
45      *      FIRST(x).
46      *
47      * 2) given lhs->...y... where y is a nonterminal symbol preceded by any
48      *      number (including 0) of nullable nonterminal symbols or actions, add
49      *      FIRST(y) to FIRST(lhs).
50      */
51
52     PRODUCTION *prod;          /* Pointer to one production side */
53     SYMBOL    **y;             /* Pointer to one element of production */
54     static SET *set = NULL;    /* Scratch-space set. */
55     int i;
56
57     if( !ISNONTERM(lhs) )      /* Ignore entries for terminal symbols */
58         return;
59

```



Listing 4.27. continued...

```

60     if( !set )                      /* Do this once. The set isn't free()'d */
61         set = newset();
62
63     ASSIGN( set, lhs->first );
64
65     for( prod = lhs->productions; prod ; prod = prod->next )
66     {
67         if( prod->non_acts <= 0 )          /* No non-action symbols      */
68         {                                     /* add epsilon to first set */
69             ADD( set, EPSILON );
70             continue;
71         }
72
73         for( y = prod->rhs, i = prod->rhs_len; --i >= 0; y++ )
74         {
75             if( ISACT( *y ) )                  /* pretend acts don't exist */
76             continue;
77
78             if( ISTERM( *y ) )
79                 ADD( set, (*y)->val );
80
81             else if( *y )                  /* it's a nonterminal */
82                 UNION( set, (*y)->first );
83
84             if( !NULLABLE( *y ) )          /* it's not a nullable nonterminal */
85                 break;
86         }
87     }
88
89     if( !IS_EQUIVALENT(set, lhs->first) )
90     {
91         ASSIGN( lhs->first, set );
92         Did_something = 1;
93     }
94 }
95
96 /*-----*/
97 PUBLIC int      first_rhs( dest, rhs, len )
98 SET           *dest;                      /* Target set            */
99 SYMBOL        **rhs;                      /* A right-hand side   */
100 int          len;                       /* # of objects in rhs */
101 {
102     /* Fill the destination set with FIRST(rhs) where rhs is the right-hand side
103     * of a production represented as an array of pointers to symbol-table
104     * elements. Return 1 if the entire right-hand side is nullable, otherwise
105     * return 0.
106     */
107
108     if( len <= 0 )
109     {
110         ADD( dest, EPSILON );
111         return 1;
112     }
113
114     for(; --len >= 0 ; ++rhs )
115     {
116         if( ISACT( rhs[0] ) )
117             continue;
118

```



Listing 4.27. continued...

```

119     if( ISTERM( rhs[0] ) )
120         ADD( dest, rhs[0]->val );
121     else
122         UNION( dest, rhs[0]->first );
123
124     if( !NULLABLE( rhs[0] ) )
125         break;
126 }
127
128 return( len < 0 );
129 }
```

Listing 4.28. follow.c— Find FOLLOW Sets

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/set.h>
4 #include <tools/hash.h>
5 #include <tools/compiler.h>
6 #include <tools/l.h>
7 #include "parser.h"
8 #include "llout.h"
9
10 /* FOLLOW.C           Compute FOLLOW sets for all productions in a
11 *                     symbol table. The FIRST sets must be computed
12 *                     before follow() is called.
13 */
14
15 int      follow_closure P(( SYMBOL *lhs ));          /* local */
16 int      remove_epsilon P(( SYMBOL *lhs ));          /* local */
17 void    init          P(( SYMBOL *lhs ));          /* local */
18 void    follow        P(( void ));                  /* public */
19
20 /*-----*/
21
22 PRIVATE int      Did_something;
23
24 /*-----*/
25
26 PUBLIC void follow()
27 {
28     D( int pass = 0; )
29     D( printf( "Initializing FOLLOW sets\n" ); )
30
31     ptab( Symtab, init, NULL, 0 );
32
33     do {
34         /* This loop makes several passes through the entire grammar, adding
35         * FOLLOW sets. The follow_closure() routine is called for each grammar
36         * symbol, and sets Did_something to true if it adds any elements to
37         * existing FOLLOW sets.
38     */
39
40     D( printf( "Closure pass %d\n", ++pass ); )
41     D( fprintf( stderr, "%d\n", pass ); )
42 }
```



Listing 4.28. continued...

```

43     Did_something = 0;
44     ptab( Symtab, follow_closure, NULL, 0 );
45
46 } while( Did_something );
47
48 /* This last pass is just for nicety and could probably be eliminated (may
49 * as well do it right though). Strictly speaking, FOLLOW sets shouldn't
50 * contain epsilon. Nonetheless, it was much easier to just add epsilon in
51 * the previous steps than try to filter it out there. This last pass just
52 * removes epsilon from all the FOLLOW sets.
53 */
54
55 ptab( Symtab, remove_epsilon, NULL, 0 );
56 D( printf("Follow set computation done\n"); )
57 }
58
59 /*-----*/
60
61 PRIVATE void init( lhs )
62 SYMBOL *lhs;           /* Current left-hand side */
63 {
64     /* Initialize the FOLLOW sets. This procedure adds to the initial follow set
65      * of each production those elements that won't change during the closure
66      * process. Note that in all the following cases, actions are just ignored.
67      *
68      * (1) FOLLOW(start) contains end of input ($).
69      *
70      * (2) Given s->...xY... where x is a nonterminal and Y is
71      *      a terminal, add Y to FOLLOW(x). x and Y can be separated
72      *      by any number (including 0) of nullable nonterminals or actions.
73      *
74      * (3) Given x->...xy... where x and y are both nonterminals,
75      *      add FIRST(y) to FOLLOW(x). Again, x and y can be separated
76      *      by any number of nullable nonterminals or actions.
77      */
78
79 PRODUCTION *prod;    /* Pointer to one production */
80 SYMBOL **x;          /* Pointer to one element of production */
81 SYMBOL **y;
82
83 D( printf("%s:\n", lhs->name); )
84
85 if( !ISNONTERM(lhs) )
86     return;
87
88 if( lhs == Goal_symbol )
89 {
90     D( printf( "\tAdding _EOI_ to FOLLOW(%s)\n", lhs->name ); )
91     ADD( lhs->follow, _EOI_ );
92 }
93
94 for( prod = lhs->productions; prod ; prod = prod->next )
95 {
96     for( x = prod->rhs ; *x ; x++ )
97     {
98         if( ISNONTERM(*x) )
99         {
100             for( y = x + 1 ; *y ; ++y )
101             {
102                 if( ISNONTERM(*y) )
103                 {
104                     for( z = y + 1 ; *z ; ++z )
105                     {
106                         if( ISNONTERM(*z) )
107                         {
108                             ADD( prod->follow, *z );
109                         }
110                     }
111                 }
112             }
113         }
114     }
115 }
116 }
```

Listing 4.28. continued...

```

102             if( ISACT(*y) )
103                 continue;
104
105             if( ISTERM(*y) )
106             {
107                 D( printf("\tAdding %s ", (*y)->name ); )
108                 D( printf("to FOLLOW(%s)\n", (*x)->name ); )
109
110                 ADD( (*x)->follow, (*y)->val );
111                 break;
112             }
113         else
114         {
115             UNION( (*x)->follow, (*y)->first );
116             if( !NULLABLE(*y) )
117                 break;
118             }
119         }
120     }
121 }
122 }
123 }
124
125 /*-----*/
126
127 PRIVATE follow_closure( lhs )
128 SYMBOL *lhs;           /* Current left-hand side */
129 {
130     /* Adds elements to the FOLLOW sets using the following rule:
131     *
132     * Given s->...x or s->...x... where all symbols following
133     * x are nullable nonterminals or actions, add FOLLOW(s) to FOLLOW(x).
134     */
135
136     PRODUCTION *prod;      /* Pointer to one production side */
137     SYMBOL    **x, **y; /* Pointer to one element of production */
138
139     D( printf("%s:\n", lhs->name ); )
140
141     if( ISACT(lhs) || ISTERM(lhs) )
142         return;
143
144     for( prod = lhs->productions; prod ; prod = prod->next )
145     {
146         for( x = prod->rhs + prod->rhs_len; --x >= prod->rhs ; )
147         {
148             if( ISACT(*x) )
149                 continue;
150
151             if( ISTERM(*x) )
152                 break;
153
154             if( !subset( (*x)->follow, lhs->follow ) )
155             {
156                 D( printf("\tAdding FOLLOW(%s) ", lhs->name ); )
157                 D( printf("to FOLLOW(%s)\n", (*x)->name ); )
158
159                 UNION( (*x)->follow, lhs->follow );
160                 Did_something = 1;
161             }

```



Listing 4.28. continued...

```

162             if( ! NULLABLE(*x) )
163                 break;
164             }
165         }
166     }
167 }
168 */
169 /*-----*/
170
171 PRIVATE remove_epsilon( lhs )
172 SYMBOL *lhs;
173 {
174     /* Remove epsilon from the FOLLOW sets. The presence of epsilon is a
175      * side effect of adding FIRST sets to the FOLLOW sets willy nilly.
176      */
177
178     if( ISNONTERM(lhs) )
179         REMOVE( lhs->follow, EPSILON );
180 }
```

Listing 4.29. llselect.c—Find LL(1) Select Sets

```

1 #include <stdio.h>
2
3 #include <tools/debug.h>
4 #include <tools/set.h>
5 #include <tools/hash.h>
6 #include <tools/compiler.h>
7 #include <tools/l.h>
8 #include "parser.h"
9
10 /* LLSELECT.C    Compute LL(1) select sets for all productions in a
11  *                  symbol table. The FIRST and FOLLOW sets must be computed
12  *                  before select() is called. These routines are not used
13  *                  by occs.
14 */
15
16 /*-----*/
17
18 void find_select_set    P(( SYMBOL *lhs ));        /* local */
19 void select              P(( void ));           /* public */
20
21 /*-----*/
22
23 PRIVATE void    find_select_set( lhs )
24 SYMBOL *lhs;                      /* left-hand side of production */
25 {
26     /* Find the LL(1) selection set for all productions attached to the
27      * indicated left-hand side (lhs). The first_rhs() call puts the FIRST sets
28      * for the initial symbols in prod->rhs into the select set. It returns true
29      * only if the entire right-hand side was nullable (EPSILON was an element
30      * of the FIRST set of every symbol on the right-hand side.
31      */
32
33     PRODUCTION *prod;
```

Listing 4.29. continued...

```

35     for( prod = lhs->productions; prod ; prod = prod->next )
36     {
37         if( first_rhs( prod->select, prod->rhs, prod->rhs_len ) )
38             UNION( prod->select, lhs->follow );
39
40         REMOVE( prod->select, EPSILON );
41     }
42 }
43
44 /*-----*/
45
46 PUBLIC void select( )
47 {
48     /* Compute LL(1) selection sets for all productions in the grammar */
49
50     if( Verbose )
51         printf("Finding LL(1) select sets\n");
52
53     ptab( Symtab, find_select_set, NULL, 0 );
54 }
```

Listing 4.30. *llcode.c*—Routines That Make the Parse Tables

```

1  #include <stdio.h>
2  #include <stdlib.h>
3  #include <ctype.h>
4
5  #include <tools/debug.h>
6  #include <tools/set.h>
7  #include <tools/hash.h>
8  #include <tools/compiler.h>
9  #include <tools/l.h>
10 #include "parser.h"
11
12 /*          LLCODE.C      Print the various tables needed for a llama-generated
13 *                           LL(1) parser. The generated tables are:
14 *
15 *           Yyd[][]      The parser state machine's DFA transition table. The
16 *                           horizontal axis is input symbol and the vertical axis is
17 *                           top-of-stack symbol. Only nonterminal TOS symbols are in
18 *                           the table. The table contains the production number of a
19 *                           production to apply or -1 if this is an error
20 *                           transition.
21 *
22 *           YydN[]       (N is 1-3 decimal digits). Used for compressed tables
23 *                           only. Holds the compressed rows.
24 *
25 *           Yy_pushtab[]  Indexed by production number, evaluates to a list of
26 *                           objects to push on the stack when that production is
27 *                           replaced. The YypDD arrays are the lists of objects
28 *                           and Yy_pushtab is an array of pointers to those lists.
29 *
30 *           Yy_snonterm[] For debugging, indexed by nonterminal, evaluates to the
31 *                           name of the nonterminal.
32 *           Yy_sact[]     Same but for the acts.
33 *           Yy_synch[]   Array of synchronization tokens for error recovery.
34 *           Yy_act()      Subroutine containing the actions.
```

Listing 4.30. continued...

```

35   *
36   * Yy_stok[] is make in stok.c. For the most part, the numbers in these tables
37   * are the same as those in the symbol table. The exceptions are the token
38   * values, which are shifted down so that the smallest token has the value 1
39   * (0 is used for EOI).
40   */
41
42 #define DTRAN           "Yyd"          /* Name of DFA transition table */
43                               /* array in the PARSE_FILE. */
44
45 /*-----*/
46
47 extern void tables()           /* public */
48
49 static void fill_row()         /* local */
50 static void make_pushtab();
51 static void make_yy_pushtab();
52 static void make_yy_dtran();
53 static void make_yy_synch();
54 static void make_yy_snonterm();
55 static void make_yy_sact();
56 static void make_acts();
57 static void make_yy_act();
58
59 /*-----*/
60 PRIVATE int *Dtran;           /* Internal representation of the parse table.
61                                     * Initialization in make_yy_dtran() assumes
62                                     * that it is an int [it calls memiset()]. */
63
64 /*-----*/
65
66 PUBLIC void tables()
67 {
68     /* Print the various tables needed by the parser. */
69
70     make_yy_pushtab();
71     make_yy_dtran();
72     make_yy_act();
73     make_yy_synch();
74     make_yy_stok();
75     make_token_file();
76
77     output( "\n#ifndef YYDEBUG\n" );
78
79     make_yy_snonterm();
80     make_yy_sact();
81
82     output( "\n#endif\n" );
83 }
84
85 /*-----*/
86 /* fill_row()
87 *
88 * Make one row of the parser's DFA transition table.
89 * Column 0 is used for the EOI condition; other columns
90 * are indexed by nonterminal (with the number normalized
91 * for the smallest nonterminal). That is, the terminal
92 * values in the symbol table are shifted downwards so that
93 * the smallest terminal value is 1 rather than MINTERM.
94 * The row indexes are adjusted in the same way (so that

```

Listing 4.30. continued...

```

95      * row 0 is used for MINNONTERM).
96      *
97      * Note that the code assumes that Dtran consists of byte-size
98      * cells.
99      */
100
101 PRIVATE void      fill_row( lhs, junk )           /* Current left-hand side */
102 SYMBOL  *lhs;
103 void      *junk;
104 {
105     PRODUCTION  *prod;
106     int          *row;
107     int          i;
108     int          rowsize;
109
110    if( !ISNONTERM(lhs) )
111        return;
112
113    rowsize = USED_TERMS + 1;
114    row    = Dtran + ((i = ADJ_VAL(lhs->val)) * rowsize );
115
116    for( prod = lhs->productions; prod ; prod = prod->next )
117    {
118        next_member( NULL );
119        while( (i = next_member(prod->select)) >= 0 )
120        {
121            if( row[i] == -1 )
122                row[i] = prod->num;
123            else
124                error(NONFATAL, "Grammar not LL(1), select-set conflict in " \
125                                "<%s>, line %d\n", lhs->name, lhs->lineno );
126        }
127    }
128 }
129
130 /*-----*/
131
132 PRIVATE void      make_pushtab( lhs, junk )
133 SYMBOL  *lhs;
134 void      *junk;
135 {
136     /* Make the pushtab. The right-hand sides are output in reverse order
137      * (to make the pushing easier) by stacking them and then printing
138      * items off the stack.
139     */
140
141     register int  i ;
142     PRODUCTION  *prod ;
143     SYMBOL      **sym ;
144     SYMBOL      *stack[ MAXRHS ], **sp;
145
146     sp = &stack[-1];
147     for( prod = lhs->productions ; prod ; prod = prod->next )
148     {
149         output( "YYPRIVATE int Yyp%02d[]={ ", prod->num );
150
151         for( sym = prod->rhs, i = prod->rhs_len ; --i >= 0 ; )
152             *++sp = *sym++;
153

```



Listing 4.30. continued...

```

154         for( ; INBOUNDS(stack,sp) ; output("%d, ", (*sp--)>val) )
155             ;
156         output( "0 );\n", prod->rhs[0] );
157     }
158 }
159 */
160
161 /*-----*/
162
163 PRIVATE void    make_yy_pushtab()
164 {
165     /* Print out yy_pushtab. */
166
167     register int      i;
168     register int      maxprod = Num_productions - 1;
169     static   char    *text[] =
170     {
171         "The YypNN arrays hold the right-hand sides of the productions, listed",
172         "back to front (so that they are pushed in reverse order), NN is the",
173         "production number (to be found in the symbol-table listing output",
174         "with a -s command-line switch).",
175         "",
176         "Yy_pushtab[] is indexed by production number and points to the",
177         "appropriate right-hand side (YypNN) array.",
178         NULL
179     };
180
181     comment( Output, text );
182     ptab   ( Symtab, make_pushtab, NULL, 0 );
183
184     output( "\nYYPRIVATE int  *Yy_pushtab[] =\n{\n" );
185
186     for( i = 0; i < maxprod; i++ )
187         output( "\tYyp%02d,\n", i );
188
189     output( "\tYyp%02d\n};\n", maxprod );
190 }
191
192 /*-----*/
193
194 PRIVATE void    make_yy_dtran( )
195 {
196     /* Print the DFA transition table. */
197
198     register int      i;
199     int            nterms, nnonterm;
200     static char    *text[] =
201     {
202         "Yyd[][] is the DFA transition table for the parser. It is indexed",
203         "as follows:",
204         "",
205         "          Input symbol",
206         "          +-----+",
207         "          L | Production number |",
208         "          H |          or YYF   |",
209         "          S |                  |",
210         "          +-----+",
211         "",
212         "The production number is used as an index into Yy_pushtab, which",
213         "looks like this:", →

```

Listing 4.30. continued...

```

214      "",  

215      "      Yy_pushtab      YypDD:,  

216      "      +-----+      +-----+",
217      "      |  *---|---->|      |",
218      "      +-----+      +-----+",
219      "      |  *---|---->",
220      "      +-----+",  

221      "      |  *---|---->",
222      "      +-----+",  

223      "",  

224      "YypDD is the tokenized right-hand side of the production.",  

225      "Generate a symbol table listing with llama's -l command-line",  

226      "switch to get both production numbers and the meanings of the",  

227      "YypDD string contents.",  

228      NULL  

229  };  

230  

231  nterms    = USED_TERMS + 1;           /* +1 for EOI */  

232  nnonterm = USED_NONTERMS;  

233  

234  i = nterms * nnonterm;             /* Number of cells in array */  

235  

236  if( !(Dtran = (int *) malloc(i * sizeof(*Dtran)) ) ) /* number of bytes */  

237      ferr("Out of memory\n");  

238  

239  memiset( Dtran, -1, i );           /* Initialize Dtran to all failures */  

240  ptab( Symtab, fill_row, NULL, 0 ); /* and fill nonfailure transitions. */  

241  

242  comment( Output , text );        /* Print header comment */  

243  

244  if( Uncompressed )  

245  {  

246      fprintf( Output,"YYPRIVATE YY_TTYPE  %s[ %d ][ %d ] =\n",
247                  DTRAN, nnonterm, nterms );  

248  

249      print_array( Output, Dtran, nnonterm, nterms );
250      defnext     ( Output, DTRAN );  

251  

252      if( Verbose )
253          printf("%d bytes required for tables\n", i * sizeof(YY_TTYPE) );
254  }
255  else
256  {
257      i = pairs( Output, Dtran, nnonterm, nterms, DTRAN, Threshold, 1);
258      pnext( Output, DTRAN );  

259  

260      if( Verbose )
261          printf("%d bytes required for compressed tables\n",
262                  (i * sizeof(YY_TTYPE)) + (nnonterm * sizeof(YY_TTYPE*)));
263  }
264  

265  output("\n\n");
266 }  

267  

268 /*-----*/  

269  

270 PRIVATE void      make_yy_synch()  

271 {  

272     int  mem ; /* current member of synch set   */  

273     int  i;      /* number of members in set      */

```



Listing 4.30. continued...

```

274
275     static char *text[] =
276     {
277         "Yy_synch[] is an array of synchronization tokens. When an error is",
278         "detected, stack items are popped until one of the tokens in this",
279         "array is encountered. The input is then read until the same item is",
280         "found. Then parsing continues.",
281         NULL
282     };
283
284     comment( Output, text );
285     output( "YYPRIVATE int Yy_synch[] =\n{\n" );
286
287     i = 0;
288     for( next_member(NULL); (mem = next_member(Synch)) >= 0; )
289     {
290         output( "\t%s,\n", Terms[mem]->name );
291         ++i;
292     }
293
294     if( i == 0 ) /* No members in synch set */
295         output( "\t_EOI_\n" );
296
297     output( "\t-1\n};\n" );
298     next_member( NULL );
299 }
300
301 /*-----*/
302
303 PRIVATE void make_yy_snonterm()
304 {
305     register int i;
306
307     static char *the_comment[] =
308     {
309         "Yy_snonterm[] is used only for debugging. It is indexed by the",
310         "tokenized left-hand side (as used for a row index in Yyd[]) and",
311         "evaluates to a string naming that left-hand side.",
312         NULL
313     };
314
315     comment( Output, the_comment );
316
317     output( "char *Yy_snonterm[] =\n{\n" );
318
319     for( i = MINNONTERM; i <= Cur_nonterm; i++ )
320     {
321         if( Terms[i] )
322             output( "\t/* %d */ \"%s\"", i, Terms[i]->name );
323
324         if( i != Cur_nonterm )
325             outc( ',' );
326
327         outc( '\n' );
328     }
329
330     output( "};\n\n" );
331 }
332

```



Listing 4.30. continued...

```

333 /*-----*/
334
335 PRIVATE void      make_yy_sact()
336 {
337     /* This subroutine generates the subroutine that implements the actions. */
338
339     register int  i;
340     static char *the_comment[] =
341     {
342         "Yy_sact[] is also used only for debugging. It is indexed by the",
343         "internal value used for an action symbol and evaluates to a string",
344         "naming that token symbol.",
345         NULL
346     };
347
348     comment( Output, the_comment );
349     output("char *Yy_sact[] =\n{\n\t");
350
351     if( Cur_act < MINACT )
352         output("NULL      /* There are no actions */");
353     else
354         for( i = MINACT; i <= Cur_act; i++ )
355         {
356             output( "\"%d\" \"%c\"", i - MINACT, i < Cur_act ? ',' : ' ' );
357             if( i % 10 == 9 )
358                 output("\n\t");
359         }
360     output("\n};\n");
361
362 }
363
364 /*-----*/
365
366 PRIVATE void      make_acts( lhs, junk )           /* Current left-hand side */
367 SYMBOL    *lhs;
368 void      *junk;
369 {
370     /* This subroutine is called indirectly from yy_act, through the subroutine
371      * ptab(). It prints the text associated with one of the acts.
372      */
373
374     char *p;
375     int  num;          /* The N in $N */
376     char *do_dollar();
377     char fname[80], *fp;
378     int  i;
379
380     if( !lhs->string )
381         return;
382
383     output( "        case %d:\n", lhs->val );
384
385     if( No_lines )
386         output( "\t\t" );
387     else
388         output( "#line %d \"%s\"\n\t\t", lhs->lineno, Input_file_name );
389
390     for( p = lhs->string ; *p ; )
391     {

```



Listing 4.30. continued...

```

392         if( *p == '\r' )
393             continue;
394
395         if( *p != '$' )
396             output( "%c", *p++ );
397         else
398         {
399             /* Skip the attribute reference. The if statement handles $$ the
400             * else clause handles the two forms: $N and $-N, where N is a
401             * decimal number. When you hit the do_dollar call (in the output()
402             * call), "num" holds the number associated with N, or DOLLAR_DOLLAR
403             * in the case of $$.
404             */
405
406         if( *++p != '<' )
407             *fname = '\0';
408         else
409         {
410             ++p;           /* skip the < */
411             fp = fname;
412
413             for(i=sizeof(fname); --i>0   && *p && *p != '>'; *fp++ = *p++ )
414             ;
415             *fp = '\0';
416
417             if( *p == '>' )
418                 ++p;
419         }
420
421         if( *p == '$' )
422         {
423             num = DOLLAR_DOLLAR;
424             ++p;
425         }
426         else
427         {
428             num = atoi( p );
429             if( *p == '-' )
430                 ++p ;
431             while( isdigit(*p) )
432                 ++p ;
433         }
434
435         output( "%s", do_dollar(num, 0, 0, NULL, fname) );
436     }
437 }
438
439     output( "\n"           break;\n" );
440 }
441
442 /*----- */
443
444 PRIVATE void    make_yy_act()
445 {
446     /* Print all the acts inside a subroutine called yy_act() */
447
448     static char *comment_text[] =
449     {
450         "Yy_act() is the action subroutine. It is passed the tokenized value",
451         "of an action and executes the corresponding code.",

```



Listing 4.30. continued...

```

452         NULL
453     };
454
455     static char *top[] =
456     {
457         "YYPRIVATE int yy_act( actnum )",
458         "{",
459         "    /* The actions. Returns 0 normally but a nonzero error code can",
460         "     * be returned if one of the acts causes the parser to terminate",
461         "     * abnormally.",
462         "    */",
463         "",
464         "    switch( actnum )",
465         "    {",
466         NULL
467     };
468     static char *bottom[] =
469     {
470         "        default: printf(\"INTERNAL ERROR: \"\
471                     \"Illegal act number (%s)\\n\", actnum);",
472         "            break;",
473         "        }",
474         "        return 0;",
475     }",
476     NULL
477 };
478
479     comment( Output, comment_text );
480     printv ( Output, top );
481     ptab   ( Symtab, make_acts, NULL, 0 );
482     printv ( Output, bottom );
483 }
```

Listing 4.31. *lldriver.c*—Routines That Make Driver Subroutines

```

1 #include <stdio.h>
2 #include <stdlib.h>
3
4 #include <tools/debug.h>
5 #include <tools/set.h>
6 #include <tools/hash.h>
7 #include <tools/compiler.h>
8 #include <tools/l.h>
9 #include "parser.h"
10
11 /-----*/
12
13 extern void file_header      P(( void ));           /* public */
14 extern void code_header       P(( void ));           /* public */
15 extern void driver           P(( void ));           /* public */
16
17 /-----*/
18
19 PRIVATE FILE    *Driver_file = stderr ;
20
```



Listing 4.31. continued...

```

21  /*-----
22   * Routines in this file are llama specific. There's a different version
23   * of all these routines in yydriver.c.
24  *-----
25  */
26
27 PUBLIC void file_header()
28 {
29     /* This header is printed at the top of the output file, before
30      * the definitions section is processed. Various #defines that
31      * you might want to modify are put here.
32     */
33
34     output( "#include <stdio.h>\n\n" );
35
36     if( Public )
37         output( "#define YYPRIVATE\n" );
38
39     if( Debug )
40         output( "#define YYDEBUG\n" );
41
42
43     output( "\n/*-----*/\n\n" );
44
45 }
46
47 /*-----
48
49 PUBLIC void code_header()
50 {
51     /* This header is output after the definitions section is processed,
52      * but before any tables or the driver is processed.
53     */
54
55     output( "\n\n/*-----*/\n\n" );
56     output( "#include \"%s\"\n\n", TOKEN_FILE );
57     output( "#define YY_MINTERM    1\n" );
58     output( "#define YY_MAXTERM    %d\n", Cur_term );
59     output( "#define YY_MINNONTERM %d\n", MINNONTERM );
60     output( "#define YY_MAXNONTERM %d\n", Cur_nonterm );
61     output( "#define YY_START_STATE %d\n", MINNONTERM );
62     output( "#define YY_MINACT    %d\n", MINACT );
63     output( "\n" );
64
65     if( !( Driver_file = driver_1(Output, !No_lines, Template) ) )
66         error( NONFATAL, "%s not found--output file won't compile\n", Template );
67 }
68
69 /*-----
70
71 PUBLIC void driver()
72 {
73     /* Print out the actual parser by copying the file llama.par
74      * to the output file.
75     */
76
77     driver_2( Output, !No_lines );
78     fclose( Driver_file );
79 }
```

Listing 4.32. *stok.c*—Routines to Make *yyout.handy_yy_stok[]*

```

1 #include <stdio.h>
2
3 #include <tools/debug.h>
4 #include <tools/set.h>
5 #include <tools/hash.h>
6 #include <tools/compiler.h>
7 #include <tools/l.h>
8 #include "parser.h"
9
10 /*-----*/
11
12 extern void make_yy_stok P(( void )); /* public */
13 extern void make_token_file P(( void ));
14
15 /*-----*/
16
17 PUBLIC void make_yy_stok()
18 {
19     /* This subroutine generates the Yy_stok[] array that's
20      * indexed by token value and evaluates to a string
21      * representing the token name. Token values are adjusted
22      * so that the smallest token value is 1 (0 is reserved
23      * for end of input).
24      */
25
26     register int i;
27
28     static char *the_comment[] =
29     {
30         "Yy_stok[] is used for debugging and error messages. It is indexed",
31         "by the internal value used for a token (as used for a column index in",
32         "the transition matrix) and evaluates to a string naming that token.",
33         NULL
34     };
35
36     comment( Output, the_comment );
37
38     output( "char *Yy_stok[] =\n{\\n" );
39     output( "\\t/* 0 */ \\\"_EOI_\",\\n" );
40
41     for( i = MINTERM; i <= Cur_term; i++ )
42     {
43         output( "\\t/* %3d */ \\\"%s\\\"", (i-MINTERM)+1, Terms[i]->name );
44         if( i != Cur_term )
45             outc( ',' );
46
47         if( (i & 0x1) == 0 || i == Cur_term )           /* Newline for every */
48             outc( '\\n' );                            /* other element */
49     }
50
51     output( "};\\n\\n" );
52 }
53
54 /*-----*/
55
56 PUBLIC void make_token_file()
57 {

```



Listing 4.32. continued...

```

58     /* This subroutine generates the yytokens.h file. Tokens have
59      * the same value as in make_yy_stok(). A special token
60      * named _EOI_ (with a value of 0) is also generated.
61      */
62
63     FILE *tokfile;
64     int i;
65
66     if( !(tokfile = fopen( TOKEN_FILE , "w") ) )
67         error( FATAL, "Can't open %s\n", TOKEN_FILE );
68
69     D( else if( Verbose ) )
70     D(     printf("Generating %s\n", TOKEN_FILE ); )
71
72     fprintf( tokfile, "#define _EOI_      0\n");
73
74     for( i = MINTERM; i <= Cur_term; i++ )
75         fprintf( tokfile, "#define %-10s %d\n",
76                         Terms[i]->name, (i-MINTERM)+1 );
77
78     fclose( tokfile );
79 }
```

Listing 4.33. main.c— Command-Line Parsing and main()

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <stdarg.h>
4 #include <signal.h>
5 #include <malloc.h>
6 #include <errno.h>
7 #include <time.h>
8 #include <sys/types.h>
9 #include <sys/timeb.h>
10 #include <process.h>
11
12 #include "date.h"
13
14 #include <tools/debug.h>
15 #include <tools/set.h>
16 #include <tools/hash.h>
17 #include <tools/compiler.h>
18 #include <tools/l.h>
19
20 #ifdef LLAMA
21 #    define ALLOCATE
22 #    include "parser.h"
23 #    undef ALLOCATE
24 #else
25 #    define ALLOCATE
26 #    include "parser.h"
27 #    undef ALLOCATE
28#endif
29
30 PRIVATE int Warn_exit = 0; /* Set to 1 if -W on command line */
31 PRIVATE int Num_warnings = 0; /* Total warnings printed */
32 PRIVATE char *Output_fname = "?????"; /* Name of the output file */
33 PRIVATE FILE *Doc_file = NULL; /* Error log & machine description */
```

Listing 4.33. continued...

```

34
35 #define VERBOSE(str)  if(Verbose){ printf( str ":\n" ); }else
36
37 /*-----*/
38
39 void onintr      P(( void ));                                /* local */
40 int  parse_args  P(( int argc,  char **argv )); 
41 int  do_file     P(( void )); 
42 void symbols    P(( void )); 
43 void statistics P(( FILE *fp )); 
44 void tail       P(( void )); 
45
46 int  main        P(( int argc,  char **argv ));             /* public */
47 void output      P(( char *fmt, ... )); 
48 void lerror      P(( int fatal, char *fmt, ... )); 
49 void error       P(( int fatal, char *fmt, ... )); 
50 char *open_errmsg P(( void )); 
51 char *do_dollar  P(( int num, int rhs_size, int lineno, PRODUCTION *prod, \
52                      char *fname)); 
53
54 /*-----*/
55 * There are two versions of the following subroutines--used in do_file(),
56 * depending on whether this is llama or occs.
57 * The occs versions is discussed in the next Chapter.
58 *
59 * subroutine:          llama version in:          occs version in:
60 *
61 * file_header()        lldriver.c                yydriver.c
62 * code_header()        lldriver.c                yydriver.c
63 * driver()            lldriver.c                yydriver.c
64 * tables()            llcode.c                 yycode.c
65 * patch()              -----                  YYPATCH.C
66 * select()             llselect.c               -----
67 * do_dollar()          lldollar.c               yydollar.c
68 *
69 * Also, several part of this file are compiled only for llama, others only for
70 * occs. The llama-specific parts are arguments to LL() macros, the occs-
71 * specific parts are in OX() macros. We'll look at what the occs-specific
72 * parts actually do in the next Chapter.
73 */
74
75
76 PUBLIC main( argc, argv )
77 char  **argv;
78 {
79     _amblksiz = 2048;           /* Declared in malloc.h */ 
80
81     signon();                 /* Print sign on message */ 
82     signal( SIGINT, onintr ); /* Close output files on Ctrl-Break. */ 
83     parse_args( argc, argv );
84
85     if( Debug && !Symbols )
86         Symbols = 1;
87
88     OX( if( Make_parser ) ) 
89     OX( { ) 
90     OX(     if( Verbose == 1 ) ) 
91     OX(     { ) 
92     OX(         if( !(Doc_file = fopen( DOC_FILE, "w" ) ) ) 
93     OX(             ferr( "Can't open log file %s\n", DOC_FILE ); ) 

```



Listing 4.33. continued...

```

94      OX(      )                                )
95      OX(    else if( Verbose > 1 )          )
96      OX(        Doc_file = stderr;           )
97      OX(      )                                )
98
99      if( Use_stdout )
100     {
101         Output_fname = "/dev/tty" ;
102         Output      = stdout;
103     }
104   else
105   {
106     OX( Output_fname = !Make_parser ? ACT_FILE : PARSE_FILE ; )
107     LL( Output_fname = PARSE_FILE;                      )
108
109     if( (Output = fopen( Output_fname, "w")) == NULL )
110       error( FATAL, "Can't open output file %s: %s\n",
111               Output_fname, open_errmsg() );
112   }
113
114   if( ( yynerrs = do_file()) == 0)      /* Do all the work */
115   {
116     if( Symbols )
117       symbols();                         /* Print the symbol table */
118
119     statistics( stdout );                /* And any closing-up statistics. */
120
121     if( Verbose && Doc_file )
122     {
123       OX( statistics( Doc_file );      )
124     }
125   }
126   else
127   {
128     if( Output != stdout )
129     {
130       fclose( Output );
131       if( unlink( Output_fname ) == -1 )
132         perror( Output_fname );
133     }
134   }
135
136 /* Exit with the number of hard errors (or, if -W was specified, the sum
137 * of the hard errors and warnings) as an exit status. Doc_file and Output
138 * are closed implicitly by exit().
139 */
140
141   exit( yynerrs + (Warn_exit ? Num_warnings : 0) );
142 }
143
144 /*-----*/
145
146 PRIVATE void onintr()                  /* SIGABRT (Ctrl-Break, ^C) Handler */
147 {
148   if( Output != stdout )              /* Destroy parse file so that a */
149   {                                    /* subsequent compile will fail */
150     fclose( Output );
151     unlink( Output_fname );
152   }
153

```


Listing 4.33. continued...

```

214      OX(           Template      = ACT_TEMPL;          )
215      OX(           break;           )                  )
216
217      case 'D': Debug       = 1;      break;
218      case 'g': Public      = 1;      break;
219      LL( case 'f': Uncompressed = 1; break;          )
220      case 'l': No_lines    = 1;      break;
221      case 'm': Template    = p + 1; goto out;
222      OX( case 'p': Make_actions = 0; break;          )
223      case 's': Symbols     = 1;      break;
224      case 'S': Symbols     = 2;      break;
225      case 't': Use_stdout  = 1;      break;
226      case 'T': Make_yyoutab = 1;    break;
227      case 'v': Verbose     = 1;      break;
228      case 'V': Verbose     = 2;      break;
229      case 'w': No_warnings = 1;      break;
230      case 'W': Warn_exit   = 1;      break;
231      LL( case 'c': Threshold = atoi( ++p );          )
232      LL(           while( *p && isdigit( p[1] ) )          )
233      LL(           ++p;          )                  )
234      LL(           break;          )                  )
235      default :               fprintf(stderr, "<-%c>: illegal argument\n", *p);
236                      printv (stderr, usage_msg );
237                      exit( EXIT_ILLEGAL_ARG );
238
239      }
240
241  out: ;
242
243  }
244
245  if( Verbose > 1 )
246      Use_stdout = 1;
247
248  if( argc <= 0 )           /* Input from standard input */
249      No_lines = 1;
250
251  else if( argc > 1 )
252  {
253      fprintf( stderr, "Too many arguments.\n" );
254      printv ( stderr, usage_msg );
255      exit ( EXIT_TOO_MANY );
256  }
257  else                      /* argc == 1, input from file */
258  {
259      if( ii_newfile( Input_file_name = *argv ) < 0 )
260      {
261          sprintf( name_buf, "%0.70s.%s", *argv, DEF_EXT );
262
263          if( ii_newfile( Input_file_name = name_buf ) < 0 )
264              error( FATAL, "Can't open input file %s or %s: %s\n",
265                      *argv, name_buf, open_errmsg() );
266      }
267  }
268
269
270 /*-----*/
271

```



Listing 4.33. continued...

```

272 PRIVATE int      do_file()
273 {
274     /* Process the input file. Return the number of errors. */
275
276     struct timeb start_time, end_time ;
277     long      time;
278
279     ftime( &start_time );    /* Initialize times now so that the difference */
280     end_time = start_time; /* between times will be 0 if we don't build the */
281     /* tables. Note that I'm using structure assign- */
282     /* ment here. */
283
284     init_acts();           /* Initialize the action code. */
285     file_header();         /* Output #defines that you might want to change */
286
287     VERBOSE( "parsing" );
288
289     nows();                /* Make lex ignore white space until ws() is called */
290     yyparse();              /* Parse the entire input file */
291
292     if( !( yynerrs || problems() ) )        /* If no problems in the input file */
293     {
294         VERBOSE( "analyzing grammar" );
295
296         first();                  /* Find FIRST sets, */
297         LL( follow(); )          /* FOLLOW sets, */
298         LL( select(); )          /* and ll(1) select sets if this is llama */
299
300         code_header();           /* Print various #defines to output file */
301         OX( patch(); )           /* Patch up the grammar (if this is occs) */
302         /* and output the actions. */
303
304         ftime( &start_time );
305
306         if( Make_parser )
307         {
308             VERBOSE( "making tables" );
309             tables();                /* generate the tables, */
310         }
311
312         ftime( &end_time );
313         VERBOSE( "copying driver" );
314
315         driver();                 /* the parser, */
316
317         if( Make_actions )
318             tail();                /* and the tail end of the source file. */
319     }
320
321     if( Verbose )
322     {
323         time = ( end_time.time * 1000 ) + end_time.millitm ;
324         time -= ( start_time.time * 1000 ) + start_time.millitm ;
325         printf( "time required to make tables: %ld.%-03ld seconds\n",
326                         (time/1000), (time%1000));
327     }
328
329     return yynerrs;
330 }
```

Listing 4.34. main.c— Error, Output, and Statistics Routines

```

331 PRIVATE void symbols( void )                                /* Print the symbol table */
332 {
333     FILE *fd;
334
335     if( !(fd = fopen( SYM_FILE, "w")) ) 
336         perror( SYM_FILE );
337     else
338     {
339         print_symbols ( fd );
340         fclose      ( fd );
341     }
342 }
343
344 /*-----*/
345
346 PRIVATE void statistics( fp )
347 FILE *fp;
348 {
349     /* Print various statistics
350     */
351
352     int conflicts;           /* Number of parse-table conflicts */
353
354     if( Verbose )
355     {
356         fprintf (fp, "\n");
357         fprintf (fp, "%4d/%-4d terminals\n",      USED_TERMS,      NUMTERMS    );
358         fprintf (fp, "%4d/%-4d nonterminals\n",    USED_NONTERMS,  NUMNONTERMS);
359         fprintf (fp, "%4d/%-4d productions\n",      Num_productions, MAXPROD   );
360     LL(         fprintf (fp, "%4d      actions\n",      (Cur_act - MINACT) +1    ) );
361     OX(         lr_stats(fp) );
362     }
363
364     LL( conflicts = 0;          )
365     OX( conflicts = lr_conflicts(fp);   )
366
367     if( fp == stdout )
368         fp = stderr;
369
370     if( Num_warnings - conflicts > 0 )
371         fprintf(fp, "%4d      warnings\n", Num_warnings - conflicts);
372
373     if(yynerrs)
374         fprintf(fp, "%4d      hard errors\n", yynerrs   );
375 }
376
377 /*-----*/
378
379 PUBLIC void output( fmt, ... )
380 char *fmt;
381 {
382     /* Works like printf(), but writes to the output file. See also: the
383      * outc() macro in parser.h
384      */
385
386     va_list args;
387     va_start( args, fmt );
388     vfprintf( Output, fmt, args );
389 }
```



Listing 4.34. continued...

```
390
391 /*-----*/
392
393 PUBLIC void document( fmt, ... )
394 char *fmt;
395 {
396     /* Works like printf() but writes to yyout.doc (provided that the file
397      * is being created.
398      */
399
400     va_list args;
401
402     if( Doc_file )
403     {
404         va_start( args, fmt );
405         vfprintf( Doc_file, fmt, args );
406     }
407 }
408
409 document_to( fp )
410 FILE *fp;
411 {
412     /* Change document-file output to the indicated stream, or to previous
413      * stream if fp=NULL.
414      */
415
416     static FILE *oldfp;
417
418     if( fp )
419     {
420         oldfp = Doc_file;
421         Doc_file = fp;
422     }
423     else
424         Doc_file = oldfp;
425 }
426
427 /*-----*/
428
429 PUBLIC void error( fatal, fmt, ... )
430 char *fmt;
431 {
432     /* This error-processing routine automatically generates a line number for
433      * the error. If "fatal" is true, exit() is called.
434      */
435
436     va_list args;
437     char *type;
438     extern int yylineno;
439
440     if( fatal == WARNING )
441     {
442         ++Num_warnings;
443         if( No_warnings )
444             return;
445         type = "WARNING: ";
446     }
447     else
448     {
449         type = "ERROR: ";
```



Listing 4.34. continued...

```

450         ++yyerrs;
451     }
452
453     va_start( args,    fmt);
454     fprintf ( stdout, "%s %s (%s, line %d): ", PROG_NAME, type,
455                           Input_file_name, yylineno );
456     vfprintf( stdout, fmt, args );
457
458     OX( if( Verbose && Doc_file ) )
459     OX( {
460         OX(     fprintf ( Doc_file, "%s (line %d) ", type, yylineno );
461         OX(     vfprintf( Doc_file, fmt, args );
462         OX(   )
463     }
464
465
466     if( fatal == FATAL )
467         exit( EXIT_OTHER );
468 }
469
470 PUBLIC void error( fatal, fmt, ... )
471 char *fmt;
472 {
473     /* This error routine works like lerror() except that no line number is
474      * generated. The global error count is still modified, however.
475      */
476
477     va_list args;
478     char *type;
479
480     if( fatal == WARNING )
481     {
482         ++Num_warnings;
483         if( No_warnings )
484             return;
485         type = "WARNING: ";
486     }
487     else
488     {
489         type = "ERROR: ";
490         ++yyerrs;
491     }
492
493     va_start ( args,    fmt  );
494     fprintf ( stdout, type );
495     vfprintf ( stdout, fmt, args );
496
497     OX( if( Verbose && Doc_file ) )
498     OX( {
499         OX(     fprintf ( Doc_file, type );
500         OX(     vfprintf ( Doc_file, fmt, args );
501         OX(   )
502
503     if( fatal == FATAL )
504         exit( EXIT_OTHER );
505 }
506
507 PUBLIC char *open_errmsg()
508 {

```



Listing 4.34. continued...

```

509     /* Return an error message that makes sense for a bad open */
510
511     extern errno;
512     switch( errno )
513     {
514         case EACCES:           return "File is read only or a directory";
515         case EXIST:            return "File already exists";
516         case EMFILE:           return "Too many open files";
517         case ENOENT:           return "File not found";
518         default:              return "Reason unknown";
519     }
520 }
521
522 /*-----*/
523
524 PRIVATE void tail()
525 {
526     /* Copy the remainder of input file to standard output. Yyparse will have
527      * terminated with the input pointer just past the %%. Attribute mapping
528      * ($$ to Yyval, $N to a stack reference, etc.) is done by the do_dollar()
529      * call.
530      *
531      * On entry, the parser will have read one token too far, so the first
532      * thing to do is print the current line number and lexeme.
533      */
534
535     extern int yylineno;          /* LeX generated */
536     extern char *yytext;          /* LeX generated */
537     int c, i, sign;
538     char fname[80], *p; /* field name in $<...>n */
539
540     output( "%s", yytext);        /* Output newline following %% */
541
542     if( !No_lines )
543         output( "\n#line %d \"%s\"\n", yylineno, Input_file_name );
544
545     ii_unterm();                  /* Lex will have terminated yytext */
546
547     while( (c = ii_advance()) != 0 )
548     {
549         if( c == -1 )
550         {
551             ii_flush(1);
552             continue;
553         }
554         else if( c == '$' )
555         {
556             ii_mark_start();
557
558             if( (c = ii_advance()) != '<' )
559                 *fname = '\0';
560             else                                /* extract name in $<foo>1 */
561             {
562                 p = fname;
563                 for(i=sizeof(fname); --i > 0 && (c=ii_advance()) != '>');
564                     *p++ = c;
565                     *p++ = '\0';
566             }

```



Listing 4.34. continued...

```

567             if( c == '>' )           /* truncate name if necessary */
568                 c = ii_advance();
569             }
570
571             if( c == '$' )
572                 output( do_dollar( DOLLAR_DOLLAR, -1, 0, NULL, fname ) );
573             else
574             {
575
576                 if( c != '-' )
577                     sign = 1 ;
578                 else
579                 {
580                     sign = -1;
581                     c = ii_advance();
582                 }
583
584                 for( i = 0; isdigit(c); c = ii_advance() )
585                     i = (i * 10) + (c - '0');
586
587                 ii_pushback(1);
588                 output( do_dollar(i * sign, -1, ii_lineno(), NULL, fname ) );
589             }
590         }
591         else if( c != '\r' )
592             outc( c );
593     }
594 }
```

Listing 4.35. ll\$ollar.c—\$ Attribute Processing (LLama Version).

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/set.h>
4 #include <tools/hash.h>
5 #include "parser.h"
6
7 PUBLIC char *do_dollar( num, rhs_size, lineno, prod, field )
8 int      num;                                /* the N in $N */
9 int      rhs_size;                           /* not used */
10 int     lineno;                            /* not used */
11 PRODUCTION *prod;                          /* not used */
12 char    *field;                            /* not used */
13 {
14     static char buf[32];
15
16     if( num == DOLLAR_DOLLAR )
17         return "Yy_vsp->left";
18     else
19     {
20         sprintf( buf, "(Yy_vsp[%d].right)", num );    /* assuming that num */
21         return buf;                                     /* has < 16 digits */
22     }
23 }
```

Listing 4.36. *signon.c*—Print Sign-on Message

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/hash.h>
4 #include <tools/set.h>
5 #include "parser.h"
6
7 PUBLIC void signon()
8 {
9     /* Print the sign-on message. Note that since the console is opened
10      * explicitly, the message is printed even if both stdout and stderr are
11      * redirected. I'm using the ANSI __TIME__ and __DATE__ macros to get the
12      * time and date of compilation.
13      */
14
15     FILE *screen;
16
17     UX( if( !(screen = fopen("/dev/tty", "w")) ) )
18     MS( if( !(screen = fopen("con:", "w")) ) )
19         screen = stderr;
20
21     LL( fprintf(screen, "LLAMA 1.0 [\" __DATE__ \" \" __TIME__ \"]\n" ); )
22     OX( fprintf(screen, "OCCS 1.0 [\" __DATE__ \" \" __TIME__ \"]\n" ); )
23
24     fprintf( screen, "(C) __DATE__ ", Allen I. Holub. All rights reserved.\n");
25     if( screen != stderr )
26         fclose(screen);
27 }

```

4.11 Exercises

- 4.1. A standard, recursive, in-order, binary-tree traversal routine looks like this:

```

typedef struct node
{
    int key;
    struct node *left, *right;
}
NODE;

traverse( root )
NODE      *root;
{
    if( !root )
        return;

    traverse( root->left );
    printf( "%d\n", root->key );
    traverse( root->right );
}

```

Write a nonrecursive binary-tree traversal function that uses a local stack to hold pointers to the parent of the current subtree rather than using the run-time stack, as is implicit in the recursive algorithm.

- 4.2. A palindrome is a word or phrase that reads the same backwards as forwards, such as the Napoleonic palindrome: “Able was I ere I saw Elba.”

- a. Write a grammar that recognizes simple palindromes made up of digits. Each digit should appear exactly twice, and the center of the palindrome should be marked with a | character. For example, the input string 1234|4321 should be recognized.
- b. Write a program that uses a push-down automata to recognize palindromes such as the foregoing.
- c. Can you modify your grammar to recognize real palindromes, such as “able was I ere I saw Elba.” If so, implement the grammar; if not, explain why.
- 4.3. Write a computer program that translates the right-linear grammars discussed in Chapter 3 into the state-machine grammars discussed in the same chapter.
- 4.4. Determine the FIRST and FOLLOW sets for all nonterminals in the first 10 productions in the grammar for C, presented in Appendix C. Is this grammar LL(1)? Why or why not?
- 4.5. Modify the following grammar so that it is not ambiguous and is LL(1):

$$\begin{array}{lcl}
 \textit{expr} & \rightarrow & -\textit{expr} \\
 & | & * \textit{expr} \\
 & | & \textit{expr} * \textit{expr} \\
 & | & \textit{expr} / \textit{expr} \\
 & | & \textit{expr} = \textit{expr} \\
 & | & \textit{expr} + \textit{expr} \\
 & | & \textit{expr} - \textit{expr} \\
 & | & (\textit{expr})
 \end{array}$$

Operator precedence and associativity should be according to the following table (higher lines in the table are higher precedence operators).

Associativity	Operator	Description
left to right	()	parentheses for grouping (highest precedence)
left to right	- *	unary minus and pointer dereference
left to right	* /	multiplication, division
left to right	+ -	addition, subtraction
right to left	=	assignment (lowest precedence)

- 4.6. Write an LL(1) grammar for the state-machine-description language presented as an exercise in the last chapter, and implement a table-driven, top-down parser for that language.
- 4.7. Is the following grammar LL(1)? Why or why not? If it isn't, convert it to LL(1).

$$\begin{array}{lcl}
 s & \rightarrow & \textit{symbol stuff} \\
 & | & \textit{TOKEN TOKEN stuff} \\
 & | & \textit{star_list stuff} \\
 \textit{star_list} & \rightarrow & \textit{star_list STAR} \\
 & | & \epsilon \\
 \textit{symbol} & \rightarrow & \textbf{TOKEN} \\
 \textit{stuff} & \rightarrow & \textbf{ANOTHER_TOKEN}
 \end{array}$$

- 4.8. Design and implement a desk-calculator program using **LLama** and **LEX**. The program should take process an input file that contains various semicolon-terminated statements. Support C syntax and the following C operators:

```
+ - / % * ^ ~ >> <<
+= -= /= %= *= ^= ~= >>= <<=
++ --
()
```

The `-` is both unary and binary minus. The `*` is a multiply. Your calculator should support a reasonable number of variables (say, 128) with arbitrary names. A variable is declared implicitly the first time it's used. In addition to the foregoing operators, your calculator should support the following pseudo subroutine:

```
printf ("` \I{format} `", ...);
```

The `printf()` pseudo subroutine should work like the normal `printf()`, except that it needs to support only the `%d`, `%x`, `%o`, and `%f` conversions. It does not have to support field widths, left justification, and so forth.

- 4.9. Modify the calculator built in the previous exercise so that it generates the assembly language necessary to evaluate the expression, rather than actually evaluating the expression at run time.
- 4.10. Modify **LLama** so that it automatically makes the following transformations in the input grammar. Productions of the form:

```
s : symbol stuff
| symbol other stuff
;
```

should be transformed to:

```
s : symbol s'
;
s': stuff
| other stuff
;
```

The transformations should be applied recursively for as long as the initial symbols on the right-hand side are the same. For example,

```
s : symbol other
| symbol other stuff
;
```

should become:

```
s : symbol s'
;
s': other
| other stuff
;
```

should become:

```
s   : symbol s'
;
s'  : other s''
;
s'' : stuff
| /* empty */
;
```

- 4.11. Modify the interactive debugging environment in `yydebug.c` so that you can look at the actual code associated with an action. You can do this by storing all action code in a file rather than in internal buffers and keeping a pointer to the file position for the start of the code in the `occs` symbol table (rather than a string pointer). Normally, this intermediate file would be destroyed by **LLama**, but the

file should remain in place if *-D* is specified on the command line. To make this file available from the debugger, **LLama** will need to output the source code for an array indexed by production number which evaluates to the correct file position. Finally, add two commands to *yydebug.c*: The first prompts for a production number and displays the associated code in the stack window. The second causes the code to be printed in the comments window every time a reduction occurs.

- 4.12. Several arrays output by **LLama** are used only for debugging, and these arrays take up unnecessary space in the executable image. Modify **LLama** and *yydebug.c* so that this information is stored in binary in an auxiliary file that is read by the debugger at run time. Don't read the entire file into memory—use *lseek()* to find the required strings and transfer them directly to the screen.
- 4.13. The action code and the debugging-environment subroutines in *yydebug.c* could be overlays—code in an action doesn't call debugger subroutines and vice versa. Arrange for overlays to be used for the debugger to save space in the executable image.

Bottom-Up Parsing

The first part of this chapter looks at bottom-up parsing in general—at how the bottom-up parse process works, how it can be controlled by a state machine, and how the state machine itself, and the associated tables, are generated. The second part of the chapter looks at a practical implementation of the foregoing: The internals of `occs`, a `yacc`-like program that builds a bottom-up parser from a grammar, is discussed in depth, as is the `occs` output file that implements the parser. There is no extended example of how `occs` is used because Chapter Six, in which code generation is discussed, provides such an example.

Bottom-up parsers such as the ones created by `yacc` and `occs` work with a class of grammars called *LALR(1)* grammars, which are special cases of the more inclusive *LR(1)* class. As with the *LL(1)* notation, the *(1)* in *LR(1)* means that one token of lookahead is required to resolve conflicts in the parse. Similarly, the first *L* means that input is read left to right. The *R* means that a *rightmost* derivation is used to construct the parse tree—the rightmost nonterminal is always expanded instead of expanding the leftmost nonterminal, as was done in the previous chapter. The *LA* in *LALR(1)* stands for “look ahead”—information about certain lookahead characters is actually built into *LALR(1)* parse tables. I’ll explain the exact characteristics of these grammars when I discuss how bottom-up parse tables are created, below. All *LL(1)* grammars are also *LR(1)* grammars, but not the other way around.

Most *LR* grammars cannot be parsed with recursive-descent parsers, and *LR* and *LALR* grammars tend to be more useful than *LL* grammars because they are easier to write—there aren’t so many restrictions placed on them. This grammatical flexibility is a real advantage when writing a compiler. *LR* parsers have many disadvantages, however. It is difficult to write an *LR* parser unless you have a tool, such as `occs` or `yacc`, to help you construct the tables used by the parser. Also, the tables are larger than *LL* parse tables and the parser tends to be slower. *LR* error recovery is difficult to do—error recovery in `yacc`-generated parsers is notoriously bad.

LALR(1) grammars.

Advantages and disadvantages of *LR* parsers.

5.1 How Bottom-Up Parsing Works*

Bottom-up parsers use push-down automata to build the parse tree from the bottom up rather than from the top down. Instead of starting at the root, as is the case with a top-down parser, a bottom-up parser starts with the leaves and, when it's collected enough leaves, it connects them together with a common root. The process is best illustrated with an example. I'll use the following expression grammar:

0.	s	\rightarrow	e
1.	e	\rightarrow	$e + t$
2.			t
3.	t	\rightarrow	$t * f$
4.			f
5.	f	\rightarrow	(e)
6.			NUM

Goal symbol.

Bottom-up parsing with a PDA, basic algorithm.

Reduce.

Shift.

Accept.

Shift=push, reduce=pop.

The nonterminals are e (for expression), t (for term), and f (for factor). The terminals are $+$, $*$, $($, $)$, and **NUM**. (**NUM** is a number—a collection of ASCII characters in the range '0' to '9'). Multiplication is of higher precedence than addition: since the star is further down in the grammar, it will be lower on the parse tree than the plus unless parentheses force it to go otherwise. The special symbol s is called the *goal symbol*. A bottom-up grammar must have one unique goal symbol, and only one production can have that goal symbol on its left-hand side. Note that this grammar can't be parsed from the top down because it's left recursive.

A bottom-up parser is a push-down automaton; it needs a state machine to drive it and a stack to remember the current state. For now, we'll ignore the mechanics of the state machine and just watch the stack. The parse process works as follows:

- (1) If the top few items on the stack form the right-hand side of a production, pop these items and replace them with the equivalent left-hand side. This procedure (popping the number of elements on a right-hand side and pushing the left-hand side) is known as a *reduce* operation.
- (2) Otherwise, push the current input symbol onto the stack and advance the input. This operation is known as a *shift*.
- (3) If the previous operation was a reduce that resulted in the goal symbol at the top of stack, *accept* the input—the parse is complete. Otherwise, go to 1.

Don't confuse shifting and reducing with pushing and popping. A *shift* always pushes a token, but it also advances the input. A *reduce* does zero or more pops followed by a push. (There are no pops on ϵ productions.) I'll illustrate the process with a parse of $1 * (2+3)$. Initially, the stack (on the left, below) is empty and none of the input characters (on the right, below) have been read:

 $1 * (2 + 3)$

Since there's no legitimate right-hand side on the stack, apply rule two: shift a **NUM** and advance past it in the input.

NUM $* (2 + 3)$

Note that the input character 1 has been translated to a **NUM** token as part of the input process. The right-hand side of Production 6 is now at top of stack, so reduce by $f \rightarrow \text{NUM}$. Pop the **NUM** and replace it with an f :

 $* (2 + 3)$

f $* (2 + 3)$

The previous operation has put another right-hand side on the stack, this time for Production 4 ($t \rightarrow f$); reduce by popping the f and pushing a t to replace it:

	* (2 + 3)
t	* (2 + 3)

The t at top of stack also forms a right-hand side, but there's a problem here. You could apply Production 2 ($e \rightarrow t$), popping the t and replacing it with an e , but you don't want to do this because there's no production that has an e followed by a * (the next input symbol) on its right-hand side. There is, however, a production with a t followed by a * (Production 3: $t \rightarrow t*f$). So, by looking ahead at the next input symbol, the conflict is resolved in favor of the shift (rather than a reduction by Production 2) in the hope that another f will eventually get onto the stack to complete the right-hand side for Production 3. The stack now looks like this:

t *	(2 + 3)
-----	-----------

There's still no reasonable right-hand side on the stack, so shift the next input symbol to the stack:

t * (2 + 3)
-------	---------

and the 2 as well:

t * (NUM	+ 3)
-----------	-------

Now there is another legitimate right-hand side on the stack: the **NUM**. The parser can apply the same productions it used to process the earlier **NUM**:

t * (+ 3)
t * (f	+ 3)
t * (+ 3)
t * (t	+ 3)

The next input symbol is a plus sign. Since the grammar doesn't have any productions in which a plus follows a t , it applies Production 2 ($e \rightarrow t$) instead of shifting:

t * (+ 3)
t * (e	+ 3)

Again, since there's still input, shift the next input symbol rather than reducing by Production 1.

t * (e +	3)
-----------	-----

and the next one:

t * (e + NUM)
---------------	---

The number forms a legitimate right-hand side, so the parser processes it as before:

t * (e +)
t * (e + f)
t * (e +)
t * (e + t)

Another right-hand side has now appeared on the stack— $e+t$ is the right-hand side of Production 1—so the parser can apply this production, popping the entire right-hand side (three objects) and pushing the corresponding left-hand side:

t * (e +)
-----------	---

$t * (e$)
$t * ($)
$t * (e$)

As before, the parser defers the application of Production 0 because there's still input; it shifts the final input symbol onto the stack:

$t * (e)$

thereby creating another right-hand side—(e) is the right-hand side of Production 5, $f \rightarrow (e)$. Reduce by this production:

$t * (e$
$t * ($
$t *$
$t * f$

The three symbols now at top of stack form the right-hand side of Production 3 ($t \rightarrow t * f$); pop them and push the left-hand side:

$t *$
t
t

Reduce by Production 2:

e

Finally, since there's no more input, apply Production 0:

s

Relationship between stack and parse tree.

Since the goal symbol is at the top of stack, the machine accepts and the parse is complete. The entire process is summarized in Table 5.1. Figure 5.1 shows the parse tree as it's created by this parse. Notice that you can determine the state of the stack for any of the partially created trees by reading across the roots of the various subtrees. For example, the tops of the subtrees shown in Figure 5.1(m) are (from left to right) t , $*$, $($, e , $+$, and f , and these same symbols comprise the stack at the equivalent time in the parse. So, you can look at the stack as a mechanism for keeping track of the roots of the partially-collected subtrees.

5.2 Recursion in Bottom-Up Parsing*

List processing in a bottom-up parser.

Hitherto, left recursion has been something to avoid in a grammar, primarily because top-down parsers can't handle it. The situation is reversed in a bottom-up parser: Left recursion should always be used for those lists in which associativity is not an issue. Right recursion can be used for right-associative lists, but it should be used only when right associativity is required. The reason for this change can be seen by looking at how right- and left-recursive lists are parsed by a bottom-up parser. Consider a left-recursive grammar, such as the following list of numbers:

Table 5.1. A Bottom-Up Parse of $1*(2+3)$

Stack	Input	Next action	See Fig. 5.1
NUM	1 * (2 + 3)	Shift NUM	a
	* (2 + 3)	Reduce by $f \rightarrow \text{NUM}$	b
	* (2 + 3)		
<i>f</i>	* (2 + 3)	Reduce by $f \rightarrow t$	c
	* (2 + 3)		
<i>t</i>	* (2 + 3)	Shift *	d
	(2 + 3)	Shift (e
<i>t</i> *	2 + 3)	Shift NUM	f
	+ 3)	Reduce by $f \rightarrow \text{NUM}$	g
<i>t</i> *	+ 3)		
	+ 3)	Reduce by $t \rightarrow f$	h
<i>t</i> *	+ 3)		
	+ 3)	Reduce by $e \rightarrow t$	i
<i>t</i> *	+ 3)		
	+ 3)	Shift +	j
<i>t</i> *	+ 3)	Shift NUM	k
)	Reduce by $f \rightarrow \text{NUM}$	l
<i>t</i> *)		
)	Reduce by $t \rightarrow f$	m
<i>t</i> *)		
)	Reduce by $e \rightarrow e+t$	n
<i>t</i> *)		
)	Shift)	o
<i>t</i> *)	Reduce by $f \rightarrow (e)$	p
<i>t</i> *		Reduce by $t \rightarrow t*f$	q
<i>t</i>		Reduce by $e \rightarrow t$	r
<i>e</i>		Reduce by $s \rightarrow e$	s
<i>s</i>		Accept	t

1. $list \rightarrow list \text{ NUM}$

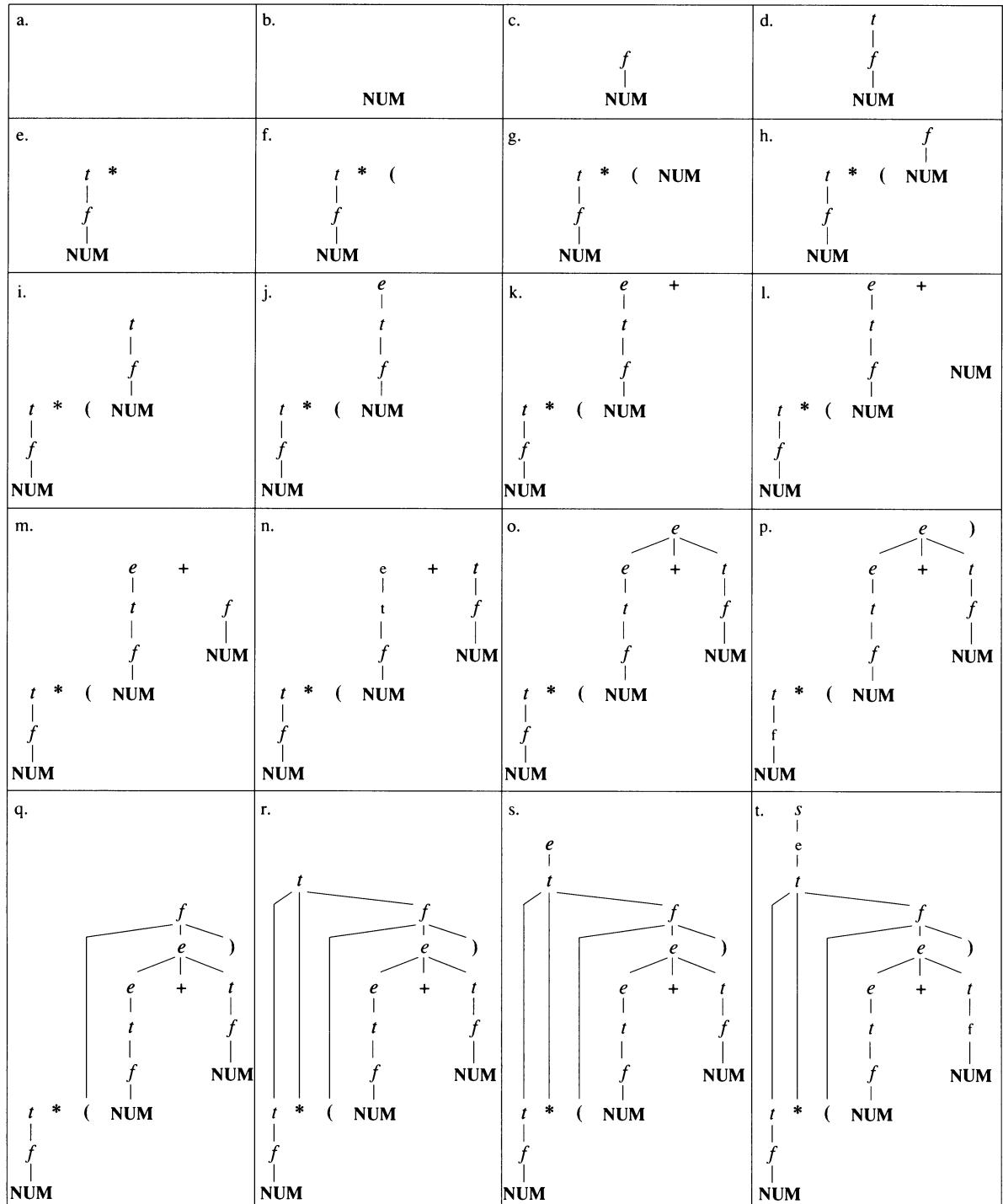
2. $| \text{ NUM}$

The input 1 2 3 generates the following bottom-up parse:

Stack	Input	Comments
—	1 2 3	
NUM	2 3	shift a NUM
<i>list</i>	2 3	reduce by $list \rightarrow \text{NUM}$
<i>list</i> NUM	3	shift a NUM
<i>list</i>	3	reduce by $list \rightarrow list \text{ NUM}$
<i>list</i> NUM		shift a NUM
<i>list</i>		reduce by $list \rightarrow list \text{ NUM}$

The stack never grows by more than two elements because a reduction occurs every time that a new **NUM** is shifted. Note that the nonrecursive production, $list \rightarrow \text{NUM}$ is applied first because the parser must get a *list* onto the stack in order to process the rest of the list elements. This is always the case in left-recursive lists such as the foregoing, and is

Nonrecursive list element reduced first.

Figure 5.1. Evolution of Parse Tree for $1 * (2 + 3)$ 

handy when doing code generation because it provides a hook for initialization actions.
(More on this in the next chapter.)

Now consider a right-recursive list, such as the following:

1. $list \rightarrow NUM\ list$
2. | NUM

Right-recursive lists use a lot of stack.

The input 1 2 3 now generates the following bottom-up parse:

Stack	Input	Comments
-	1 2 3	Shift a NUM
NUM	2 3	Shift a NUM
NUM NUM	3	Shift a NUM
NUM NUM NUM		Apply $list \rightarrow NUM$
NUM NUM list		Apply $list \rightarrow NUM$ list
NUM list		Apply $list \rightarrow NUM$ list
list		Apply $list \rightarrow NUM$ list

Here, all of the list elements are pushed onto the stack before any of them are processed via a reduction. This behavior is always the case—right-recursive productions use a lot of stack because all the list elements must be pushed before a reduction can occur, so it's worth avoiding right associativity if at all possible. Of course, you can't do this in an expression grammar because certain operators just have to associate in certain ways. Lists, however, are another matter.

5.3 Implementing the Parser as a State Machine*

I need to introduce a few terms at this point. At the risk of repeating myself, a *shift* is the process of moving a token from the input stream to the stack—push the current token and advance the input. A *reduce* is the process of replacing a right-hand side on the stack with the matching left-hand side—pop as many items as there are on the right-hand side of a production and push the left-hand side. A *viable prefix* is a partially parsed expression—an incomplete derivation. In a bottom-up parser, the symbols on the parse stack from bottom to top, are identical to the viable prefix, read left to right.

Shift, reduce, viable prefix.

A *handle* is a right-hand side on the top few items of the stack. The parser scans the stack looking for a handle, which when found, triggers a reduce operation. Put another way, a reduce can be performed when the rightmost symbols of the viable prefix form a right-hand side. Strictly speaking, a handle is a right-hand side which, when found on the stack, triggers a reduction. The parser might defer a reduction even though a right-hand side is at the top of the stack—the associativity and precedence rules built into the grammar might require that more input symbols be read before any reductions occur. In this situation, a right-hand side does not form a handle because it doesn't trigger a reduction. In most situations, however, a reduction occurs as soon as the right-hand side appears on the stack.

Handle.

Not all right-hand sides in viable prefix form handles.

To automate a bottom-up parser, some mechanism is needed for recognizing handles as they appear on the stack. Of course, you could do this by brute force—scanning the stack and comparing what you find there to a list of right-hand sides—but this method is so inefficient as to be impractical. It's better to build a state machine that keeps track of all the push and pop operations. The condition of the stack can then be inferred from a knowledge of what items have been pushed or popped. The accepting states in the machine trigger a reduce operation. Nonaccepting states represent either a shift or the push of a left-hand side that is part of a reduce operation. The state machine in Figure 5.2. serves to illustrate some of the concepts involved. This machine implements the

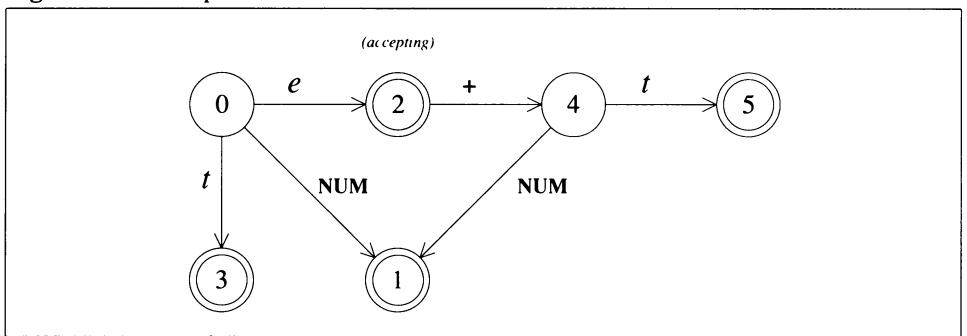
Using a state machine to recognize handles.

following grammar, a subset of the one used in the previous example.

0.	$s \rightarrow e$
1.	$e \rightarrow e + t$
2.	$ $ t
3.	$t \rightarrow \text{NUM}$

This grammar recognizes expressions composed either of single numbers or of alternating numbers and plus signs. The expressions associate from left to right. The state machine can be viewed as a sort of super syntax diagram that reflects the syntactic structure of an entire language, not just of the individual productions. It's as if you formed an NFA by combining the syntax diagrams for all the productions in the grammar.

Figure 5.2. A Simple LR State Machine



Stack remembers state transitions. Arrows indicate pushes.

How terminals and non-terminals get onto stack.

States and handles.

The automaton's stack keeps track of state transitions. The current state is the one whose number is at top of stack. The direction of the arrows in the state diagram represents pushes. (The state number associated with the next state is pushed.) Pops cause us to retrace the initial series of pushes. That is, the stack serves as a record of the various state transitions that were made to get to the current state. You can retrace those transitions backwards by popping the states. For example, starting at State 0, a zero is at the top of the stack. Shifting a **NUM** results in a transition from State 0 to State 1, and a 1 is pushed to record this transition. Reducing by $t \rightarrow \text{NUM}$ involves a pop (which causes a retrace from State 1 back to State 0) and a push of a 3 (which causes a transition from State 0 to State 3). The parser can retrace the original path through the machine because that information is on the stack—the state that you came from is at the top of the stack after the pop part of the reduction.

It's important to notice here that the only way that a terminal symbol can get onto the stack is by the push that occurs in a shift operation, and the only way that a nonterminal symbol can get onto the stack is by the push that occurs during a reduction.

The purpose of the state machine is to detect handles. In this machine, States 1, 2, 3, and 5 represent the recognition of a handle. (A reduce from State 2 also terminates the parse successfully.) Entering one of these states usually triggers a reduce operation (depending on the next input character). For example, if the machine is in State 1, the parser must have just shifted a **NUM**, so it will reduce by $t \rightarrow \text{NUM}$. If the machine is in State 3, the parser must have just reduced by a production that had a t on its left-hand side, and it will do another reduce (by $e \rightarrow t$). If the machine's in State 2, the parser reduces by $s \rightarrow e$ only if the next input character is the end-of-input marker, otherwise it will shift to State 4. State 5 tells the parser that an e , plus, and t must have been pushed so the handle $e+t$ is recognized, and it reduces by $e \rightarrow e+t$. Since the state stack itself, along with a knowledge of what those states represent, is sufficient to see what's going on, there's no reason to keep a stack of actual symbols, as I did in the previous section.

The state machine from Figure 5.2. is represented in array form in Table 5.2—columns are possible input symbols (in the *Action* section) or nonterminal symbols (in the *Goto* section). The special symbol \vdash in the *Action* table is an end-of-input marker. The rows represent possible states of the machine. The basic parsing algorithm is shown in Table 5.3.

LR parse tables: *Action* and *Goto* components.

Table 5.2. Array Representation of an LR State Machine

		Action			Goto	
		\vdash	NUM	+	e	t
Top of Stack	0	—	s1	—	2	3
	1	r3	—	r3	—	—
	2	Accept	—	s4	—	—
	3	r2	—	r2	—	—
	4	—	s1	—	—	5
	5	r1	—	r1	—	—

I'll demonstrate the parse process by parsing the expression $1+2+3$. (The entire parse is summarized in Table 5.4—I'll do the parse step by step here, however.) The algorithm starts off by pushing 0, the state number of the start state.

Table-driven LR parse,
an example.

0 1 + 2 + 3 \vdash

Action[0][NUM] holds an *s1*, so shift a 1 onto the stack entering State 1. Advance the input as part of the shift operation:

0 1 + 2 + 3 \vdash

Now, since the new top-of-stack symbol is a 1, the parser looks at Action[1][+] and finds an *r3* (reduce by production 3). The algorithm in Figure 5.2 tells the parser to execute a code-generation action and pop one item off the stack (because there's one item on the right-hand side of $t \rightarrow \text{NUM}$).

0 + 2 + 3 \vdash

The parser now has to figure the Goto transition. The left-hand side of Production 3 is a *t* and popping the 1 uncovered the 0, so the parser consults Goto[0][*t*], finding a 3, which it pushes onto the state stack:

0 3 + 2 + 3 \vdash

Spinning back up to the top of the while loop, Action[3][+] is an *r2*, a reduction by Production 2 ($e \rightarrow t$). The parser pops one item and then pushes the contents of Goto[0][*e*], State 2:

0 + 2 + 3 \vdash

0 2 + 2 + 3 \vdash

Looping again, Action[2][+] is an *s4*—shift a 4, advancing the input:

0 2 4 2 + 3 \vdash

Action[4][NUM] is an *s1* so we shift again:

0 2 4 1 + 3 \vdash

Action[1][+] is an *r3*.

0 2 4 + 3 \vdash

This time, however, the uncovered top-of-stack item is a 4, not a 0, as it was the first time the parser reduced by Production 3. Goto[4][*t*] is a 5, so the machine goes to State 5 by

Table 5.3. State-Machine-Driven LR Parser Algorithm

<pre> push(0) while(Action[TOS][input] ≠ Accept) { if(Action[TOS][input] = -) error() else if (Action[TOS][input] = sX) push(X) advance() else if (Action[TOS][input] = rX) act(X) pop(as many items as are in the RHS of Production X) push(Goto[uncovered TOS][LHS of Production X]) } accept(); } </pre>																						
<p>Definitions:</p> <table style="width: 100%; border-collapse: collapse;"> <tr> <td>Action</td> <td>Columns in the <i>Action</i> part of the state-machine array.</td> </tr> <tr> <td>accept()</td> <td>Return success status. An accept is actually a reduction by Production 0 on end of input.</td> </tr> <tr> <td>act(X)</td> <td>Perform the action associated with a reduction of Production X. This action usually generates code.</td> </tr> <tr> <td>advance()</td> <td>Discard current input symbol and get the next one.</td> </tr> <tr> <td>error()</td> <td>Print error message and do some sort of error recovery.</td> </tr> <tr> <td>Goto</td> <td>Columns in the <i>Goto</i> part of the state-machine array.</td> </tr> <tr> <td>input</td> <td>Current input symbol.</td> </tr> <tr> <td>Stack</td> <td>State stack.</td> </tr> <tr> <td>push(X)</td> <td>Push state X onto the stack.</td> </tr> <tr> <td>pop(N)</td> <td>Pop N items from the stack.</td> </tr> <tr> <td>TOS</td> <td>State number at the top of the stack.</td> </tr> </table>	Action	Columns in the <i>Action</i> part of the state-machine array.	accept()	Return success status. An accept is actually a reduction by Production 0 on end of input.	act(X)	Perform the action associated with a reduction of Production X. This action usually generates code.	advance()	Discard current input symbol and get the next one.	error()	Print error message and do some sort of error recovery.	Goto	Columns in the <i>Goto</i> part of the state-machine array.	input	Current input symbol.	Stack	State stack.	push(X)	Push state X onto the stack.	pop(N)	Pop N items from the stack.	TOS	State number at the top of the stack.
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input	Current input symbol.																					
Stack	State stack.																					
push(X)	Push state X onto the stack.																					
pop(N)	Pop N items from the stack.																					
TOS	State number at the top of the stack.																					

pushing a 5 onto the stack:

0 2 4 5 + 3 ⊢

Looping yet again, Action[5][+] is an r1, so an *e*, plus, and *t* must be on the stack, and the parser reduces by Production 1 ($e \rightarrow e + t$). Note that the current state (5) can only be reached by making transitions on an *e*, +, and *t*. Since each forward-going transition is a push, the stack must contain these three symbols. Since there are three items on the right-hand side of the production, pop three states:

0 + 3 ⊢

and then push Goto[0][*e*] to replace them:

0 2 + 3 ⊢

Next, since Action[2][+] is a s4, the parser goes to State 4 with a shift:

0 2 4 3 ⊢

Action[4][NUM] is an s1 so it shifts to State 1:

0 2 4 1 ⊢

Action[1][-] calls for a reduction by Production 3:

0 2 4	⊤
0 2 4 5	⊤
Action[5][⊤]	is a reduction by Production 1:
0 2 4	⊤
0 2	⊤
0	⊤
0 2	⊤

Action[2][⊤] is an accept action, so the parse is finished. The tree implied by the previous parse is shown in Figure 5.3.

Table 5.4. A Bottom-Up Parse of 1+2+3 Using a State Machine

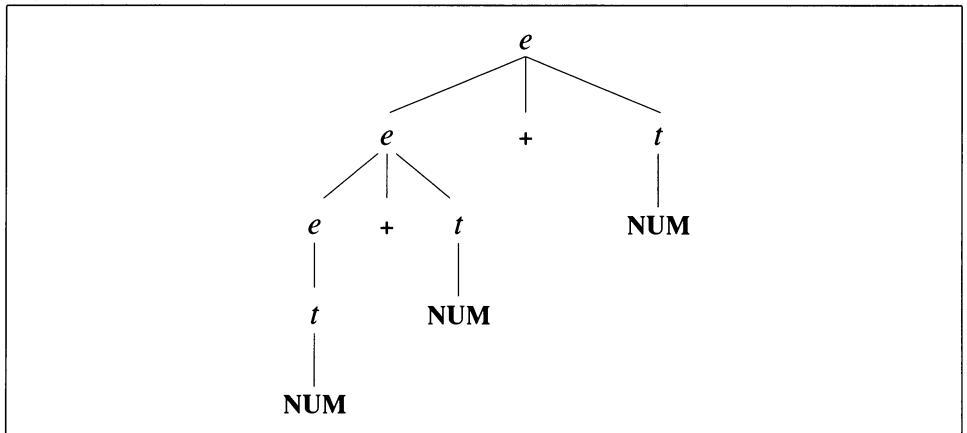
State Stack	Symbols	Input	Notes
0	\$	1 + 2 + 3 ⊤	Push start state
0 1	\$ NUM	+ 2 + 3 ⊤	Shift NUM
0	\$	+ 2 + 3 ⊤	Reduce by $t \rightarrow \text{NUM}$
0 3	\$ t	+ 2 + 3 ⊤	
0	\$	+ 2 + 3 ⊤	Reduce by $e \rightarrow t$
0 2	\$ e	+ 2 + 3 ⊤	
0 2 4	\$ e +	2 + 3 ⊤	Shift +
0 2 4 1	\$ e + NUM	+ 3 ⊤	Shift NUM
0 2 4	\$ e +	+ 3 ⊤	Reduce by $t \rightarrow \text{NUM}$
0 2 4 5	\$ e + t	+ 3 ⊤	
0 2 4	\$ e +	+ 3 ⊤	Reduce by $e \rightarrow e + t$
0 2	\$ e	+ 3 ⊤	
0	\$	+ 3 ⊤	
0 2	\$ e	+ 3 ⊤	
0 2 4	\$ e +	3 ⊤	Shift +
0 2 4 1	\$ e + NUM	⊤	Shift NUM
0 2 4	\$ e +	⊤	Reduce by $t \rightarrow \text{NUM}$
0 2 4 5	\$ e + t	⊤	
0 2 4	\$ e +	⊤	Reduce by $e \rightarrow e + t$
0 2	\$ e	⊤	
0	\$	⊤	
0 1	\$ e	⊤	Accept

There is one difference between a symbol-oriented stack, as was discussed in the previous section, and a state stack: the latter requires an initial push of the start state that wasn't required on the symbol stack. I'm using a \$ to hold the place of this extra stack item on the symbol stack. Also, note that the term *accept* is often used in two ways when PDA-based parsers are discussed. First of all, there are normal accepting states, in this case States 1, 2, 3, and 5. A transition into an accepting state signifies that a handle has been recognized. The action associated with these states is a reduce (and the associated code-generation actions, which are discussed below). You can also use the word *accept* to signify that a complete program has been parsed, however. If you look at the parser as a recognizer program, the machine as a whole *accepts* or *rejects* an input sentence. (A *sentence* in this context is a complete program, as represented by the grammar. A valid sentence in a grammar forms a complete parse tree.)

Difference between symbol and state stack.

Accept.

Reject.
Sentence.

Figure 5.3. A Parse Tree for $1+2+3$ 

5.4 Error Recovery in an LR Parser*

Stack lacks lookahead information for error recovery.

Panic-mode error recovery.

Avoid cascading error messages.

Error recovery in a bottom-up parser is notoriously difficult. The problem is that, unlike a top-down parser's stack, the bottom-up parse stack contains no information about what symbols are expected in the input. The stack tells you only what's already been seen. One effective technique, called *panic-mode* error recovery, tries to get around this difficulty by using the parse tables themselves to find legal input symbols. It works as follows:

- (0) Error recovery is triggered when an error transition is read from the parse table entry for the current lookahead and top-of-stack symbols (when there's no legal outgoing transition from a state on the current input symbol).
- (1) Remember the current condition of the stack.
- (2) Discard the state at the top of stack.
- (3) Look in the parse table and see if there's a legal transition on the current input symbol and the uncovered stack item. If so, the parser has recovered and the parse is allowed to progress, using the modified stack. If there's no legal transition, and the stack is not empty, go to (2).
- (4) If all items on the stack are discarded, restore the stack to its original condition, discard an input symbol, and go to (2).

The algorithm continues either until it can start parsing again or until the complete input file is absorbed. Table 5.5 shows what happens when the incorrect input $1++2$ is parsed using the simple LR state machine in Table 5.2 on page 345. Messages should be suppressed if a second error happens right on the tail of the first one in order to avoid cascading error messages. Nothing should be printed if an error happens within a limited number (four or five is a good choice) of parse cycles (shift or reduce operations) from the point of the previous error.

5.5 The Value Stack and Attribute Processing*

Attributes in recursive descent.

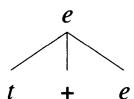
The next issue is how bottom-up parsers, such as the one generated by `occs`, handle attributes and use them to generate code. In a recursive-descent parser, attributes can be computed anywhere in the parse process, but they are only passed in two places: Inherited attributes are passed when a subroutine is called, and synthesized attributes are passed when the subroutine returns. Since a table-driven LR parser doesn't use

Table 5.5. Error Recovery for $t \rightarrow e + t$

Stack	Input	Comments
— —	1 + + 2 \vdash	Shift start state
0 \$	1 + + 2 \vdash	Shift NUM (goto 1)
0 1 \$ NUM	+ + 2 \vdash	Reduce by Production 3 ($t \rightarrow \text{NUM}$) (Return to 0, goto 3)
0 3 \$ t	+ + 2 \vdash	Reduce by Production 2 ($e \rightarrow t$) (Return to 0, goto 2)
0 2 \$ e	+ + 2 \vdash	Shift + (goto 4)
0 2 4 \$ e +	+ 2 \vdash	ERROR (no transition in table) Pop one state from stack
0 2 \$ e	+ 2 \vdash	There is a transition from 2 on + Error recovery is successful
0 2 \$ e	+ 2 \vdash	Shift + (goto 4)
0 2 4 \$ e +	2 \vdash	Shift NUM (goto 1)
0 2 4 1 \$ e + NUM	2 \vdash	Shift NUM (goto 1)
0 2 4 1 \$ e + NUM	\vdash	Reduce by Production 3 ($t \rightarrow \text{NUM}$) (Return to 4, goto 5)
0 2 4 5 \$ e + t	\vdash	Reduce by Production 1 ($e \rightarrow e+t$) (Return to 0, goto 2)
0 2 \$ e	\vdash	Accept

subroutine calls, another method is required to pass attributes and execute code-generation actions. To simplify matters, it's rarely necessary to use inherited attributes in a bottom-up parser; I'll concentrate on synthesized attributes only.

Reviewing the relationship between a bottom-up parse and the parse tree itself, a production like $e \rightarrow t+e$ generates the following tree:



In a bottom-up parse, the t is pushed onto the stack by means of a reduction by $t \rightarrow \text{NUM}$, the plus is just shifted, then there is a flurry of activity which terminates by the bottom e being pushed as part of a reduction by either $e \rightarrow e+t$ or $e \rightarrow t$. In any event, once the t , plus, and e are on the stack, a handle is recognized; all three symbols are popped, and an e (the left-hand side) is pushed. Synthesized attributes move up the parse tree—from the children to the parent, and they are passed during a reduce operation—when the handle is replaced with the left-hand side. The synthesized attribute is attached to the left-hand side that is pushed as part of a reduce. Inherited attributes, which go from a parent to a child, would be pushed as part of a shift operation. The situation is simplified by ignoring any inherited attributes, so attributes of interest are all passed during the reduce

Bottom-up parser uses synthesized attributes.

operation—when the parser moves up the tree. If you reexamine the LR parse algorithm presented in Table 5.3 on page 346, you’ll notice that code generation occurs only on a reduce operation.

Value, attribute stack.

It turns out that the easiest place to store attributes is on the parse stack itself. In practical terms, you can do this in one of two ways, either make the state stack a stack of structures, one field of which is the state number and the rest of which is for attributes, or implement a second stack that runs parallel to the state stack. This second stack is called a *value* or *attribute* stack. Every push or pop on the state stack causes a corresponding push or pop on the value stack. All code-generation actions in a bottom-up parser are executed after the parser decides to reduce, but before the parser pops anything off the stack. So, you need to arrange for any attributes that are needed by a code-generation action to be on the value stack at the appropriate time. Similarly, the code-generation action needs to be able to put attributes onto the value stack for use by subsequent actions. I’ll demonstrate with an example.

Bottom-up attribute passing, an example.

Figure 5.4 shows a parse of $1+2+3$ with both the state and value stacks shown (the value stack is on the bottom). Emitted code is printed to the right of the value stack. The attributes that appear on the value stack are the names of the temporary variables used in expression evaluation. (I’ll demonstrate how they get there shortly.) I’ve used \square to mark undefined value-stack entries.

Look, initially, at stack number two in Figure 5.4. Here, a **NUM** has been shifted onto the stack and the corresponding value-stack item is undefined. The next stage in the parse is a reduce by $t \rightarrow \text{NUM}$. Before the parser does the actual reduction, however, it executes some code. Two actions are performed: First, the $t_0=1$ instruction is generated. (The 1 is the lexeme attached to the **NUM** token that is shifted onto Stack 2 in Figure 5.4.) Next, the code that outputs the instruction leaves a note for the parser saying that the attribute t_0 should be placed on top of the value stack after the reduction is performed. When the left-hand side is pushed as the last step of a reduction, the parser must push a t_0 in the corresponding place on the value stack.

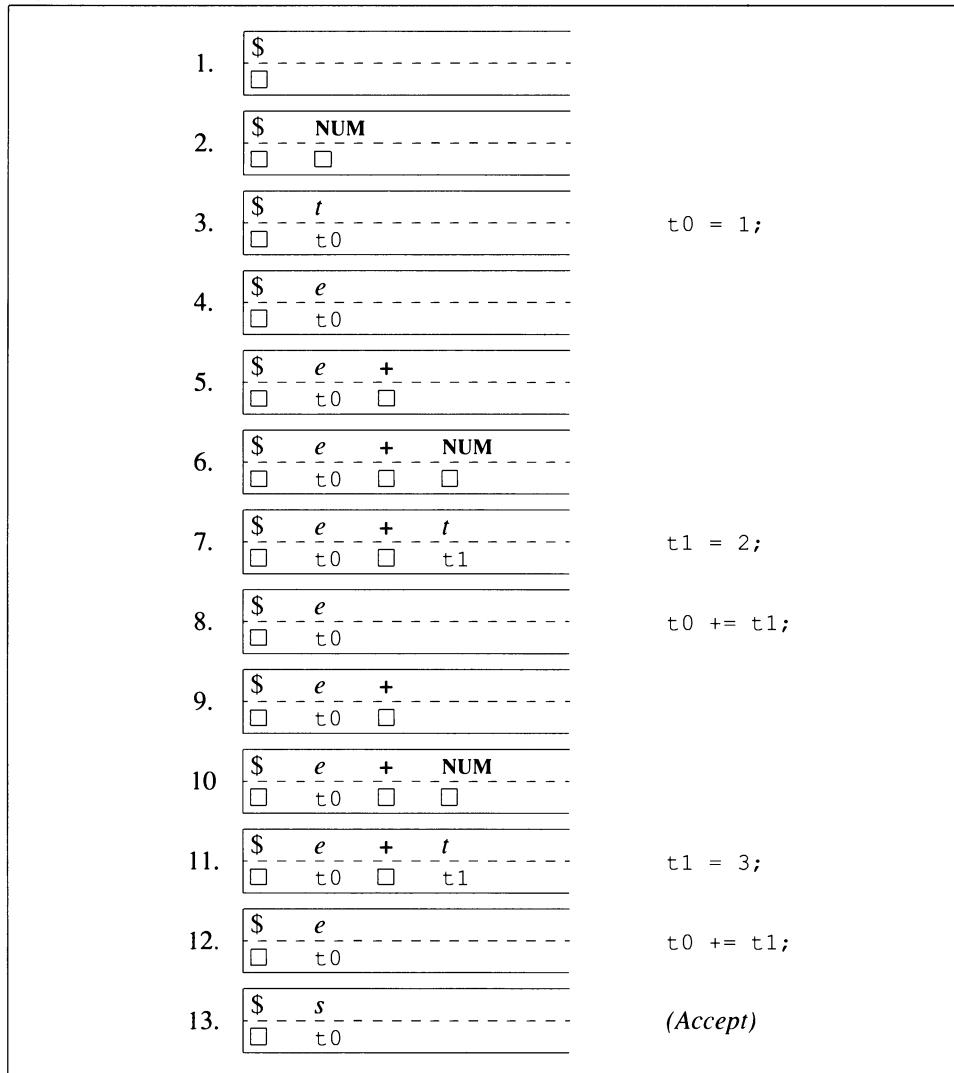
The parser now does the actual reduction, popping the **NUM** from the state stack and replacing it with a t . It modifies the value stack too, popping one item and, instead of pushing garbage, it pushes a t_0 onto the value stack. In other words, the attribute t_0 on the value stack is now attached to the t on the state stack (because it’s at the same relative position). The t has acquired the attribute t_0 .¹

The next step in the parse process is application of $e \rightarrow t$. No code is generated here. There is an associated action, however. The e must acquire the attributes formerly associated with the t . So, again, the parser leaves a note for itself saying to push t_0 onto the value stack as it pushes the new left-hand side onto the state stack.

The next transition of interest is the one from stack seven to eight. Here, the handle $e+t$ is recognized and the parser applies $e \rightarrow e+t$. Two actions are performed. A $t_0+=t_1$ instruction is generated, and the attribute t_0 is acquired by the e (is passed to the e). The important thing to notice here is that the names of the two temporary variables are attributes of e and t . The names of these variables are on the value stack. The code that generates the add instruction can discover these names by examining the cells at offsets zero and two from the top of the value stack. Moreover, the position of these attributes on the

The names of the temporaries are on the value stack.

1. Note that it’s a common, but incorrect, usage to say that the left-hand side inherits attributes from the right-hand side, probably because that’s the way that control flows through the actual code (the child’s code is executed first). I’d recommend against this usage because of the potential confusion involved. It’s better to say that a parent *acquires* an attribute (or *is passed* an attribute) that’s synthesized by the child.

Figure 5.4. The Value Stack and Code Generation

stack can be determined by examining the grammar. The handle on the stack comprises a right-hand side, so the position of the attribute on the stack corresponds to the position of the symbol on the right-hand side. The rightmost symbol is at top of stack. The e is two symbols away from the rightmost symbol in the production, so it's at offset two from the top of stack.

The synthesized attribute here is the name of the temporary variable that holds the result of the addition at run time; the t_0 that's attached to the e on stack eight is the name of the anonymous temporary that holds the result of the entire subexpression evaluation. It's just happenstance that the attribute is the same one that was attached to the e in stack seven. The e in stack eight is not the same e as the one in stack seven—the one in stack eight is the one on the left-hand side of $e \rightarrow e+t$, the e in stack seven is on the right-hand side of the same production. If the code generator had decided to emit an instruction like $t_3=t_0+t_1$, then the synthesized attribute would have been t_3 .

The next step in building a parser is to augment the earlier grammar to incorporate code-generation and attribute-passing rules. This augmentation is done in Table 5.6.

Augmented and attributed grammars for bottom-up parsing.

`$$` represents left-hand side.

The actions are performed immediately before reducing by the associated production. The special symbol `$$` is used to leave a note about what attribute to push. The value that's assigned to `$$` in an action is, in turn, attached to the left-hand side that's pushed as part of the reduce operation. (`$$` is pushed onto the value stack when the new left-hand side is pushed onto the parse stack.) The code assumes that the value stack is declared as an array of character pointers:

```
char *Value_stack[128] ;
char **Vsp = Value_stack + sizeof(Value_stack);
```

`Vsp` is the stack pointer. The stack grows towards low memory: `**Vsp=value` is a push and `value=*``Vsp++` is a pop.² Consequently, `Vsp[0]` is the item at top of stack, `Vsp[1]` is the cell immediately below that, and so forth.

Table 5.6. Expression Grammar with Actions

	Grammar	Actions
0.	$S \rightarrow e$	<code>generate("answer = %s", Vsp[0]);</code>
1.	$e \rightarrow e + t$	<code>generate("%s += %s", Vsp[2], Vsp[0]);</code> <code>free_name(Vsp[0]);</code> <code>\$\$ = Vsp[2];</code>
2.	$e \rightarrow t$	<code>\$\$ = Vsp[0];</code>
3.	$t \rightarrow \text{NUM}$	<code>name = new_name();</code> <code>generate("%s = %s", name, yytext);</code> <code>\$\$ = name</code>

Parser maintains value stack.

`generate()`,
`new_name()`,
`free_name()`.

The value stack is maintained directly by the parser. A place marker that correspond to the terminal that's pushed onto the parse stack is pushed onto the value stack during a shift. In a reduce, as many objects as are removed from the parse stack are also popped off the value stack, then `$$` is pushed to correspond with the push of the left-hand side.

The actions in Table 5.6 call three subroutines: `generate()` works like `printf()`, except that it puts a newline at the end of every line, `new_name()` returns a string that holds the name of a temporary variable, `free_name()` puts the name back into the name pool so that it can be reused by a subsequent `new_name()` call. The three routines are implemented in Listing 5.1. (Error detection has been removed for clarity.) The `Names` array is a stack of available temporary-variable names, initialized to hold eight names. `Namep` is the name-stack pointer. The `new_name()` subroutine pops a name and returns it; `free_name()` pushes the name to be freed.

-
2. Downward-growing stacks are required by ANSI C unless you want to waste one cell of the array. An ANSI pointer is assumed always to point into an array, though it can go one cell past the end of the array. It cannot point below the base address of the array, however. The following code demonstrates the problem:

```
int x[10];
p = &x[10]; /* this is valid */
p = x; /* this is valid */
--p; /* p is now undefined */
```

The problem is a real one in an 8086-family machine. If `x` is at address 1000:0000, the `--p` in a compact- or large-model program usually evaluates to 1000:fffe rather than 0fff:fffe. Consequently `p<x` is always false because 1000:fffe is larger than 1000:0000.

Listing 5.1. support.c— Support Routines for Actions in Table 5.6

```

1 #include <stdio.h>
2 #include <stdarg.h>
3
4 char *Names[] = { "t0", "t1", "t2", "t3", "t4", "t5", "t6", "t7" };
5 char **Namep = Names;
6
7 void generate( fmt, ... )
8 char *fmt;
9 {
10     va_list args;
11     va_start( args, fmt );
12     vfprintf( stdout, fmt, args );
13     fputc( '\n', stdout );
14 }
15
16 char *new_name()
17 {
18     if( Namep >= &Names[ sizeof(Names)/sizeof(*Names) ] )
19     {
20         printf("Expression too complex\n");
21         exit( 1 );
22     }
23
24     return( *Namep++ );
25 }
26
27 free_name(s)
28 char *s;
29 {
30     *--Namep = s;
31 }
```

5.5.1 A Notation for Bottom-Up Attributes*

Though attributes are actually at specific offsets from the top of the value stack, it's inconvenient always to refer to them in terms of the stack pointer. **Yacc** and **occs** both use a notation for describing attributes that is sufficiently useful to merit mention here. (This notation is used heavily in the next chapter.) As before, $\$\$$ holds the value that is pushed as part of a reduction. Other attributes are indicated with the notation $\$1$, $\$2$, and so forth, where the number represents the position of the symbol on the right-hand side, for example:

$$\begin{array}{ccccccc} \$\$ & \$1 & \$2 & \$3 & \$4 & \dots \\ s & \rightarrow & \mathcal{A} & \mathcal{B} & \mathcal{C} & \mathcal{D} & \dots \end{array}$$

In addition, **yacc** and **occs** automatically provide a default action of $\$\$ = \1 if no explicit assignment is made to $\$\$$. Don't confuse this notation with the similar notation used by **LLama**, in which the number in $\$n$ is the offset of a symbol to the right of an action. Here, n is the absolute position of the symbol on the right-hand side, not a relative position.

Table 5.7 shows the augmented expression grammar we just looked at modified to use these conventions. The actions are all enclosed in curly braces to separate them from other symbols in the production. No explicit actions are now required in Production 2 because the default $\$\$=\$1$ is used.

Bottom-up attribute notation: $\$\$$, $\$1$, $\$2$, etc.

Default action: $\$\$=\$1$.

Table 5.7. Augmented Expression Grammar with Dollar Notation

0.	$s \rightarrow e$	{ generate("answer = %s", \$1); }
1.	$e \rightarrow e + t$	{ generate("%s += %s", \$1, \$3); free_name(\$3); }
2.	$e \rightarrow t$	
3.	$t \rightarrow \text{NUM}$	{ name = new_name(); generate("%s = %s", \$\$ = name, yytext); }

5.5.2 Imbedded Actions*

Note that the actions are not really grammatical symbols in a bottom-up parser, as they were in the top-down parser—the actions are not part of the viable prefix, they are just executed at appropriate times by the parser. In particular, an action is performed only when the parser performs a reduction. Hitherto, our actions have all been to the far right of the production to signify this fact. It's convenient sometimes to put an action in the middle of a production, however, as in the following production:

```
function_body → arg_list {args_done();} compound_stmt {funct_done();}
```

In order to process this sort of production, you must modify the grammar so that the middle action can be done as part of a reduction. The easiest way to do this is to provide a dummy ϵ production as follows:

```
function_body → arg_list dummy compound_stmt {funct_done();}  
dummy → ε {args_done();}
```

Shift/reduce conflicts caused by imbedded actions.

Both **yacc** and **occs** do this transformation for you automatically, but the extra ϵ production can often create parsing problems called shift/reduce conflicts, discussed below in detail. You can often eliminate the unwanted shift/reduce conflict by inserting an imbedded action immediately to the right of a terminal symbol in the production, but it's generally better to avoid imbedded actions entirely. You could rearrange the previous grammar as follows without introducing an ϵ production:

```
function_body → formal_arg_list compound_stmt {funct_done();}  
formal_arg_list → arg_list {args_done();}
```

5.6 Creating LR Parse Tables—Theory*

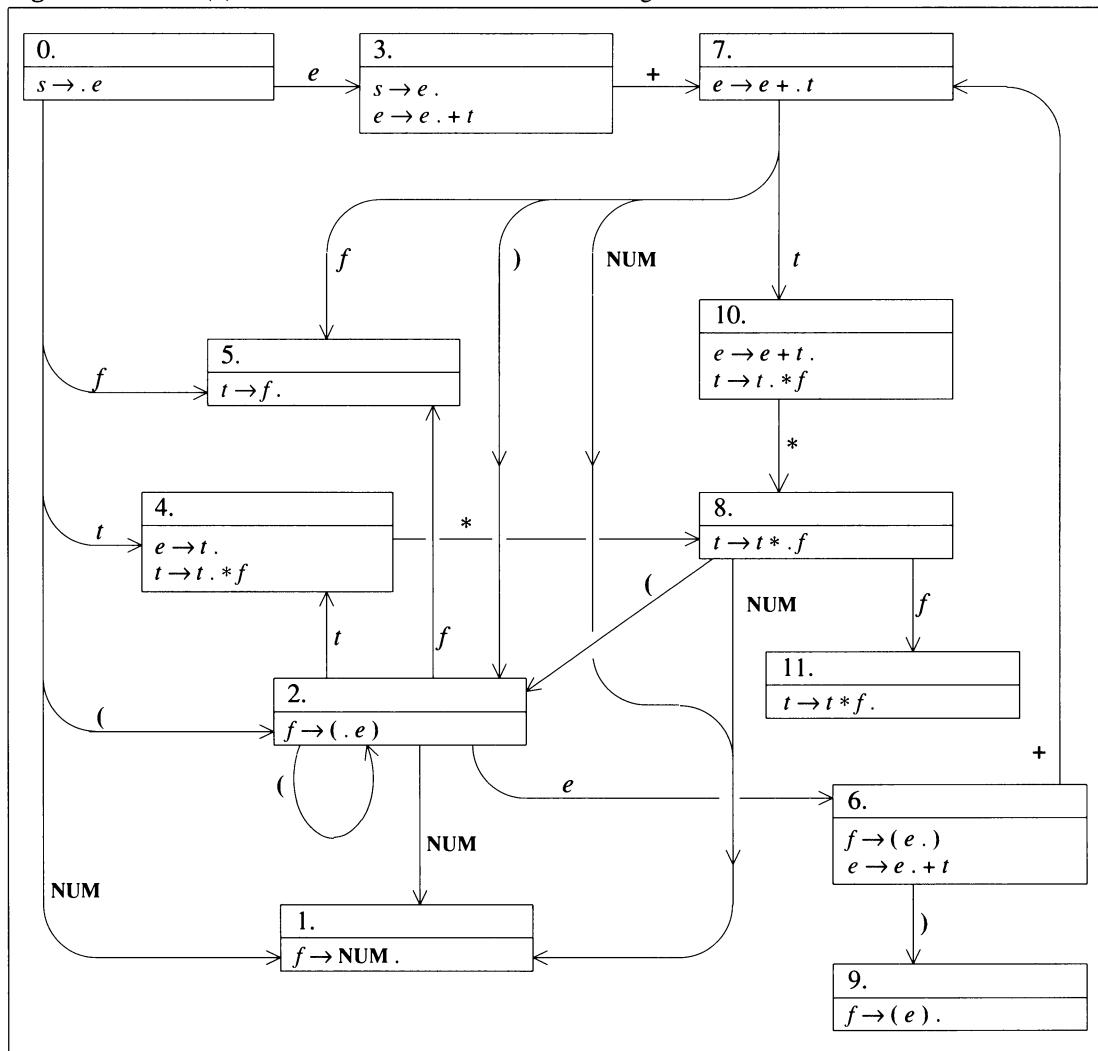
This section discusses how LR parse tables are created, and also discusses the LR-class grammars in depth.

5.6.1 LR(0) Grammars*

I'll start discussing LR-class grammars by considering two, more restrictive LR grammars and why they're too limited for most practical applications. These are LR(0) grammars—LR grammars that require no lookahead tokens—and SLR(1) grammars (*simple* LR(1) grammars). The expression grammar in Table 5.8 is used throughout the current section. The nonterminals are $+$, $*$, $($, $)$, and **NUM** (a number). Note that this grammar is not LL(1), so it can't be parsed with a top-down parser. The parse table for this grammar is represented as a graph in Figure 5.5 and in tabular form in Table 5.9.

Table 5.8. An Expression Grammar

0:	$s \rightarrow e$
1:	$e \rightarrow e + t$
2:	$e \rightarrow t$
3:	$t \rightarrow t * f$
4:	$t \rightarrow f$
5:	$f \rightarrow (e)$
6:	$f \rightarrow \text{NUM}$

Figure 5.5. LALR(1) State-Machine for Grammar in Listing 5.8

The table and graph relate to one another as follows:

- Each box is a state, and the state number is the row index in both the *action* and *goto* parts of the parse table.
- Outgoing edges labeled with nonterminal symbols represent the *goto* part of the parse table—a number representing the nonterminal is used as a column index and

How bottom-up parse table relates to state diagram.

Nonterminal edges are *goto* entries.

Table 5.9. State-Transition Table for Grammar in Table 5.8

	Shift/Reduce Table (Yy_action)						Goto Table (Yy_goto)			
	⊣	NUM	+	*	()	s	e	t	f
0	–	s1	–	–	s2	–	–	3	4	5
1	r6	–	r6	r6	–	r6	–	–	–	–
2	–	s1	–	–	s2	–	–	6	4	5
3	accept	–	s7	–	–	–	–	–	–	–
4	r2	–	r2	s8	–	r2	–	–	–	–
5	r4	–	r4	r4	–	r4	–	–	–	–
6	–	–	s7	–	–	s9	–	–	–	–
7	–	s1	–	–	s2	–	–	–	10	5
8	–	s1	–	–	s2	–	–	–	–	11
9	r5	–	r5	r5	–	r5	–	–	–	–
10	r1	–	r1	s8	–	r1	–	–	–	–
11	r3	–	r3	r3	–	r3	–	–	–	–

sN =shift to state N . rN =reduce by production N .
Error transitions are marked with –.

Terminal edges are *shift* entries.

Reduce directives.

Error transitions.

LR(0) items: Dots in productions.

Dot movement.

the table holds the next-state number. Remember that this path represents the push part of a reduction.

- All outgoing edges that are labeled with terminal symbols become shift directives in the *action* part of the parse table. The token value is used as a column index, and the parser shifts from the current to the indicated state when it encounters that token.
- Reduce directives in the parse table are found by looking at the productions in each state. A reduction happens from a state when a production has a dot at its far right. A reduction by $f \rightarrow \text{NUM}$ can happen from State 1 because State 1 has the production $f \rightarrow \text{NUM.}$ in it. (Remember, the dot to the right of the **NUM** signifies that a **NUM** has already been read and shifted onto the stack.) I'll describe how **occs** decides in which column to put the reduce directives in a moment.
- Everything else in the table is an error transition.

The productions in the bottom part of each state are *LR(0) items*; an LR(0) item is different from a production in that it is augmented with a dot, which does two things. First, it keeps track of what's on the parse stack—in a given state, everything to the left of the dot is on the stack. When a dot moves to the far right of a production, the entire right-hand side is on the stack and the parser can reduce by that production. Second, the dot shows you how much input has been read so far—everything to the right of the dot in an item for a given state has not yet been input when the parser is in that state. Consequently, symbols to the right of the dot can be used as lookahead symbols.

When you move from one state to another, the dot moves from the left of the symbol with which the edge is labeled to the right of that symbol. In the case of a terminal symbol, the dot's movement means that the parser has read the token and shifted it onto the stack. In the case of a nonterminal, the parser will have read all the input symbols that comprise that nonterminal—in terms of parse trees, the parser will have processed the entire subtree rooted at that nonterminal, and will have read all the terminal symbols that form leaves in the subtree. In other words, the dot movement signifies that the parser has pushed a nonterminal onto the stack—something that can happen only as part of a reduction. When the parser reduces, it will have processed all of the input associated with the pushed nonterminal.

Now, I'll demonstrate how particular items get into particular states. First, notice that there are two classes of states: those that have incoming edges labeled with terminal symbols, and those with incoming edges labeled with nonterminal symbols. There are no states with both kinds of incoming edge—it's either one or the other. I'll look at the two classes of states separately.

Adding items to states.

Start by considering how the parse works. From State 0, a **NUM** token causes a shift to State 1, which in turn causes the dot to move past the **NUM**. Since the parser just shifted the **NUM** to the stack, it's at the top of the stack, and the dot is to its immediate right. Since all transitions into State 5 involve a shift of a **NUM**, all items in State 5 must have a **NUM** immediately preceding the dot. This is the case with all states that have incoming edges labeled with terminal symbols. The dot follows that nonterminal in every item in the state.

You must examine the reduce operation to see what's happening with the other states. All reductions are done in two steps. First, a previous state is restored by going backwards through the machine (against the direction of the arrows) by popping as many states as there are symbols on a right-hand side of the production. Next, the parser makes a transition to a new state by following an edge labeled with the left-hand side of the production by which you're reducing. These next states comprise the goto part of the parse table.

We're interested in the second of these steps, the outgoing edges that are labeled with left-hand sides. Looking at an example, since the dot is to the far right of the production in State 5, the parser reduces from State 1 by applying $f \rightarrow \text{NUM}$. The reduction involves a single pop, which uncovers State 0, followed by a push of the f , which moves us to State 5. Looking at the items in State 5, you'll notice that the dot follows the f . Since an f must be on top of the stack in this state, the dot must be to the immediate right of the f in all items in State 5. The next reduction is by $t \rightarrow f$ and gets us, first back to State 0, and then to State 4. Since a t must now be at top of stack, all items in State 4 have a dot to the immediate right of t . Let's look at it another way. State 0 contains the item $[e \rightarrow e + t]$.³ A reduction by some production that has an e on its left-hand side causes us to return to State 0 and then go to State 3. When the parser does this reduction, it has read past all input symbols that comprise the e , and the following item appears in State 3 to signify this fact: $[e \rightarrow e . + t]$. Given the foregoing, you can create the parse table with the following procedure:

Procedure for creating bottom-up parse tables.

- (1) **Initialize:** State 0 is the start state. Initially, it contains a single LR(0) item consisting of the *start production* (the one with the goal symbol as its left-hand side), and a dot at the far left of the right-hand side. The initial set of items in this state (and all other states) is called the *kernel* or *seed* items. The set of kernel items for the start state looks like this:

0
$s \rightarrow . e$

Initialization. Start production.

Kernel, seed items.

- (2) **Close the kernel items:** The kernel items for the next states are created with a *closure* operation, which creates a set of *closure items*. Closure is performed on those kernel items that have a nonterminal symbol to the right of a dot. (I'll call this non-terminal symbol the *closure symbol*.) Add to the closure set items for those

Closure.

3. Items are usually differentiated from simple productions by surrounding them with brackets, as is done here.

productions which, when reduced, can return us to the current state.

- Initiate the closure process by adding to the closure set all productions that have the closure symbol on their left-hand side. Remember, if the dot is to the left of a nonterminal, there will be an outgoing transition from the current state (labeled with that nonterminal), and that transition represents the push part of a reduction. Consequently, productions with that nonterminal on their left-hand side could return us to the current state when a reduction by that production is performed.

The new item is formed by placing a dot at the far left of the right-hand side of the production. In the current grammar, since an e is to the right of the dot in all the kernel items, items containing all productions with an e on their left-hand side are added. The closure set looks like this so far:

0
$s \rightarrow .e$
$e \rightarrow .e + t$
$e \rightarrow .t$

- Repeat the closure process on the newly added items. If a newly added item has a nonterminal to the right of the dot, and closure has not yet been performed on that nonterminal, add to the closure set all productions that have that nonterminal as their left-hand side. In the current example, the second production has a t to the right of the dot, so I'll add all productions with t on their left-hand side, yielding the following:

0
$s \rightarrow .e$
$e \rightarrow .e + t$
$e \rightarrow .t$
$t \rightarrow .t * f$
$t \rightarrow .f$

- Repeat (b) until no more items can be added. A production with an f to the right of a dot was added in the previous step, so the closure set is extended to include all productions that have f on their left-hand sides, yielding:

0
$s \rightarrow .e$
$e \rightarrow .e + t$
$e \rightarrow .t$
$t \rightarrow .t * f$
$t \rightarrow .f$
$f \rightarrow .(e)$
$f \rightarrow .\text{NUM}$

Since none of the new productions have dots to the left of a nonterminal, the procedure is complete.

(3) Form the kernel items for the next states:

- Partition the items:** Group together all items in the current state (both kernel and closure items) that have the same symbol to the right of the dot. I'll call this symbol the *transition symbol*, and all items that share a transition symbol form a *partition*. The following picture shows the items partitioned by symbol to the right of the dot.

0.
$s \rightarrow . e$
$e \rightarrow . e + t$
$e \rightarrow . t$
$t \rightarrow . t * f$
$t \rightarrow . f$
$f \rightarrow . (e)$
$f \rightarrow . \text{NUM}$

Partitions that have dots at the far right of the productions do not form new states because they represent reductions from the current state. (There will be, at most, one of these partitions, and all ϵ productions are part of this partition.) All other partitions form the kernels of the new states.

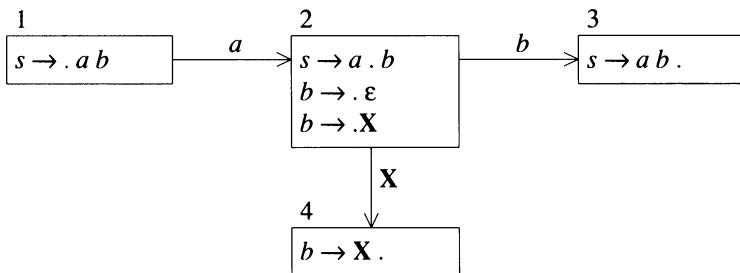
- (b) **Form the kernel items for the next states:** Modify the items in those partitions that form the new states by moving the dot one notch to the right, and:
 - (c) **Add next-state transitions:** If a state already exists that has the same set of kernel items as the modified partition, add an edge labeled with the transition symbol that goes from the current state to the existing state. Otherwise, create a new state, using the group of modified items as the kernel items. Add an edge from the current state to the new one on the transition symbol.
- (4) Repeat steps (2) and (3) on the kernel items of any new states until no more states can be added to the table.

Figure 5.6 shows a complete state machine built from our grammar using the foregoing procedure. The initial state is State 0, the first pass of the algorithm added States 1 to 5, and so on.

The issue of ϵ productions deserves further mention. ϵ productions are special only in one way. There are no outgoing edges from any state that are labeled ϵ ; rather, if an ϵ production is added to a state as part of a closure operation, the machine reduces from the current state on that production. In other words, the right-hand side of an ϵ production effectively has the dot at its far right to begin with—the ϵ may as well not be there because it represents an empty string. Consider a grammar such as the following fragment:

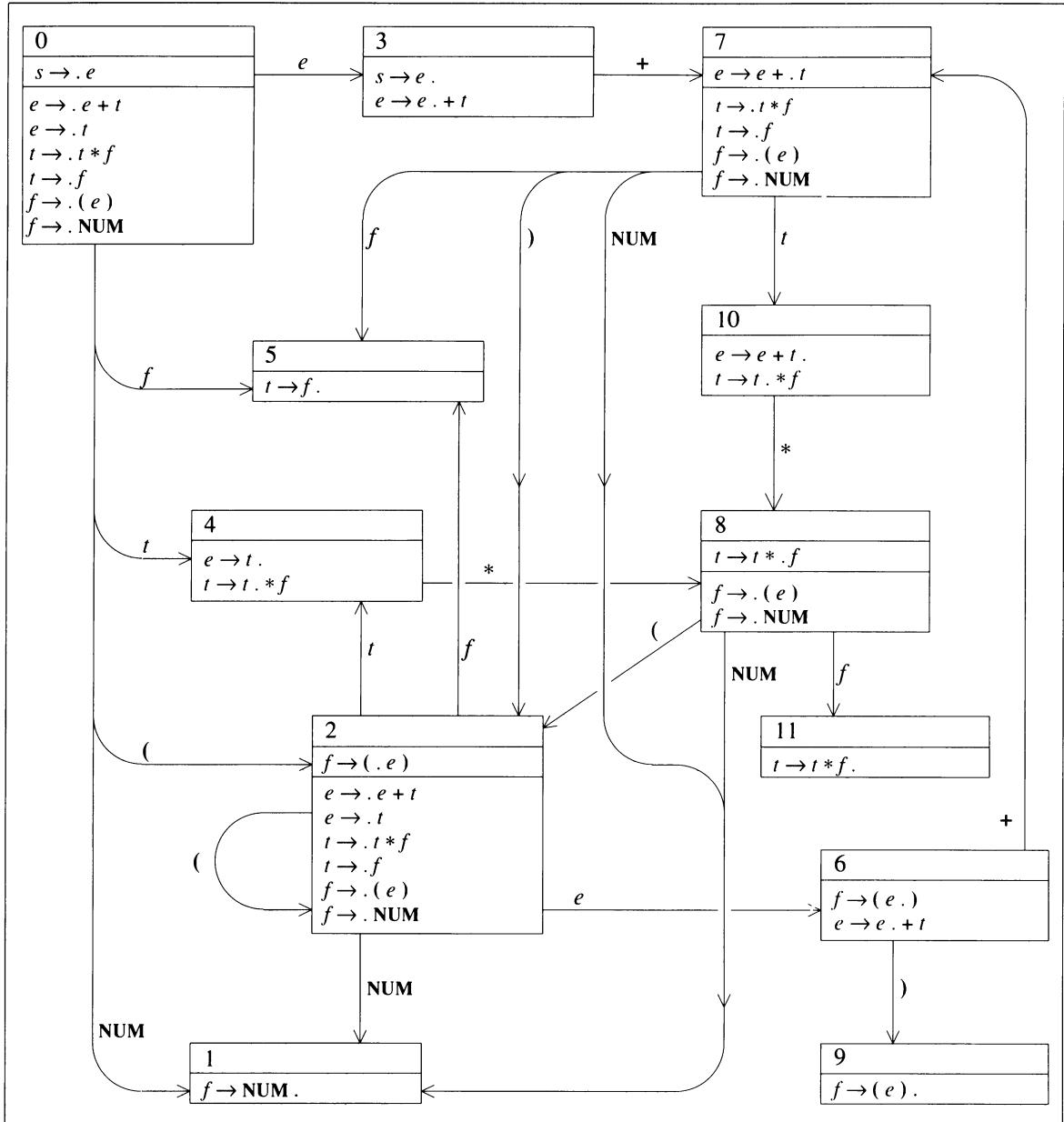
$$\begin{array}{lcl} s & \rightarrow & a b \\ b & \rightarrow & \epsilon \\ | & & X \end{array}$$

which generates the following states:



ϵ productions in bottom-up parse tables.

The item $[a \rightarrow a.b]$ adds the closure items $[b \rightarrow \epsilon.]$ and $[b \rightarrow X.]$, the latter of which appears as $[b \rightarrow X.]$ in State 4. If b goes to ϵ in State 2, a reduction is performed from that state. Since there are no symbols on the right-hand-side, you won't pop anything off the state stack, so you won't change states in the first part of the reduce. Nonetheless, since b is the left-hand side of the production, you do want to go to State 2 in the goto part of

Figure 5.6. An LR(0) State Machine

the reduce operation, so there's a transition on b from State 1.

Shift/reduce,
reduce/reduce conflicts.

Unfortunately, LR(0) state machines have problems. As I said earlier, a reduction is performed from any given state if a dot is at the far right of an item, and a terminal symbol is shifted if the dot is to its immediate left. But consider a state like State 4 in Figure 5.6, where one of the items has a dot to the far right, and the other has a dot buried in the middle of the production. Should the parser shift or reduce? This situation is called a *shift/reduce conflict*. It's also possible for a state to have two items with different productions, both of which have dots at the far right. This situation is called a conflict—I'll look at one in a moment. States that contain shift/reduce or reduce/reduce conflicts are called *inadequate states*. A grammar is LR(0) if you can use the foregoing procedures to

Inadequate states.

create a state machine that has no inadequate states. Obviously, our current grammar is not LR(0).

5.6.2 SLR(1) Grammars*

There is a simple way to resolve some shift/reduce conflicts. Remember how the FOLLOW sets were used to form LL(1) select sets in the previous chapter. If a right-hand side was nullable, you could apply that production if the lookahead symbol was in the FOLLOW set of the left-hand side, because there could be a derivation in which all elements of the right-hand side could effectively disappear, and the symbol following the left-hand side would be the next input symbol in this case.

Use FOLLOW sets to resolve conflicts.

You can use similar reasoning to resolve shift/reduce conflicts. In order to do a reduction, the next input symbol must be in the FOLLOW set of the nonterminal on the left-hand side of the production by which we're reducing. Looking at State 3 in the LR(0) machine in Figure 5.6, you can reduce by $s \rightarrow e$ if the lookahead symbol is in $\text{FOLLOW}(s)$. Similarly, you must shift if the lookahead symbol is a plus sign. As long as $\text{FOLLOW}(s)$ doesn't contain a plus sign, the shift/reduce conflict in this state can be resolved by using the next input symbol. If all inadequate states in the LR(0) machine generated from a grammar can be resolved in this way, then you have an SLR(1) grammar.

The FOLLOW sets for our current grammar are in Table 4.17. Looking at the shift/reduce conflict in State 4, $\text{FOLLOW}(e)$ doesn't contain a *, so the SLR(1) method works in this case. Similarly, in State 3, $\text{FOLLOW}(s)$ doesn't contain a +, so everything's okay. And finally, in State 10, there is an outgoing edge labeled with a *, but $\text{FOLLOW}(e)$ doesn't contain a *. Since the FOLLOW sets alone are enough to resolve the shift/reduce conflicts in all three states, this is indeed an SLR(1) grammar.

Table 5.10. FOLLOW Sets for the Expression Grammar

$\text{FOLLOW}(s)$	\equiv	$\{ \vdash \}$
$\text{FOLLOW}(e)$	\equiv	$\{ \vdash \} + \{ \}$
$\text{FOLLOW}(t)$	\equiv	$\{ \vdash \} + \{ * \}$
$\text{FOLLOW}(f)$	\equiv	$\{ \vdash \} + \{ * \}$

5.6.3 LR(1) Grammars*

Many grammars are not as tractable as the current one—it's likely that a FOLLOW set will contain symbols that also label an outgoing edge. A closer look at the machine yields an interesting fact that can be used to solve this difficulty. A nonterminal's FOLLOW set includes *all* symbols that can follow that nonterminal in every possible context. The state machine, however, is more limited. You don't really care which symbols can follow a nonterminal in every possible case; you care only about those symbols that can be in the input when you reduce by a production that has that nonterminal on its left-hand side. This set of relevant lookahead symbols is typically a subset of the complete FOLLOW set, and is called the *lookahead set*.

Some symbols in FOLLOW set are not needed.

The lookahead set associated with a nonterminal is created by following a production as it moves through the state machine, looking at those terminal symbols that can actually follow the nonterminal just after a reduction. In the case of the inadequate State 10, there are two paths from State 0 to State 10 (0–3–7–10 and 0–2–6–7–10). For the purposes of collecting lookahead sets, the only important terminals are those that follow non-terminal symbols preceded by dots. (We're looking for the sequence: dot-nonterminal-terminal.) Remember, here, that we're really figuring out which elements of the

Lookahead set.

Creating lookahead sets.

FOLLOW set tell us whether or not to reduce in a given situation. If a dot precedes a nonterminal, then that nonterminal can be pushed onto the parse stack as part of the reduction, and when the nonterminal is pushed, one of the terminals that follows it in the grammar better be in the input.

Look at State 0 in Figure 5.6. The parser can enter State 0 in the middle of a reduction by some production that has an e on its left-hand side, and it will exit State 0 by the e edge. This edge was created because a dot preceded an e in some item in that state. Looking at State 0 again, a plus can be in the input when the parser reduces by a production with e on its left-hand side because of the item:

$$[e \rightarrow \cdot e + t]$$

If the parser enters State 0 in the middle of a reduction by a production with e on its left-hand side, then a $+$ could reasonably be the next input symbol. Looked at another way, after all the tokens that comprise the e are read, $+$ could be the next input symbol in this context. There's also an implied end-of-input marker in the start production, so

$$[s \rightarrow \cdot e \vdash]$$

tells us to add \vdash to the lookahead set. The item:

$$[f \rightarrow \cdot (e)]$$

doesn't add anything to the lookahead set, even though it has an e in it, because this item isn't considered when a reduction is performed (the dot doesn't precede the e). Of the other states along the two paths from State 0 to State 10, the only other lookahead is added in State 4, because of the item:

$$[f \rightarrow (\cdot e)]$$

If the parser enters State 4 in the middle of a reduction by some production with an e on its left-hand side, then a right parenthesis is also a possible input character. When you get to State 10, the only elements of $\text{FOLLOW}(e)$ that can actually follow an e in this context are the elements that you've collected by looking at things that can actually follow an e in some state on the path from the start state to State 10. This set of lookaheads has only three elements: $+$, $)$, and \vdash . A $*$ is not part of this set, so a $*$ cannot legally follow the e in this context, and the parser doesn't have to consider the $*$ when it's deciding whether to reduce from State 10. It can safely reduce if the lookahead symbol is a $+$, $)$, or \vdash , and shift on a $*$.

LR(1) items.

Formalizing the foregoing procedure, an *LR(1) item* is an LR(0) item augmented with a lookahead symbol. It has three parts (a production, a dot, and a lookahead) and is typically represented like this: $[s \rightarrow \alpha . \beta, D]$, where D is the lookahead symbol. Note that the lookahead is part of the item; two items with the same production and dot position, but with different lookahead, are different items.

Creating LR(1) state machines.

The process of creating an LR(1) state machine differs only from that used to make an LR(0) machine only in that LR(1) items are created in the closure operation rather than LR(0) items. The initial item consists of the start production with the dot at the far left and \vdash as the lookahead character. In the grammar we've been using, it is:

$$[s \rightarrow \cdot e, \vdash]$$

An LR(1) item is created from a kernel item as follows: Given an item and a production that take the following forms:

$$\begin{array}{ll} s & \rightarrow \alpha . x \beta, C \\ x & \rightarrow \cdot \gamma \end{array}$$

(s and x are nonterminals, α , β , and γ are collections of any number of terminals and non-terminals [“any number” includes zero], and C is a terminal symbol) add: $[s \rightarrow \alpha \cdot x \beta, C]$.

$[x \rightarrow \cdot \gamma, \text{FIRST}(\beta C)]$.

to the closure set. $\text{FIRST}(\beta C)$ is computed like FIRST of a right-hand side. If β is not nullable, then $\text{FIRST}(\beta C)$ is $\text{FIRST}(\beta)$, otherwise it is the union of $\text{FIRST}(\beta)$ and C . $[x \rightarrow \cdot \gamma, \text{FIRST}(\beta C)]$.

Looking at a real situation, the start production relates to the previous formula as follows:

$$\begin{array}{c} [s \rightarrow \alpha \cdot x \beta, C] \\ \hline [s \rightarrow \cdot e, \vdash] \end{array}$$

Note that both α and β match empty strings in this situation. Using the previous formula, there are two productions that have e on the left-hand side: $e \rightarrow e+t$ and $e \rightarrow t$, so add the following to the closure set: α and β can be empty.

$$\begin{array}{c} [s \rightarrow \cdot \gamma, \text{FIRST}(\beta C)] \\ [e \rightarrow \cdot e + t, \text{FIRST}(\epsilon \vdash)] \\ [e \rightarrow \cdot t, \text{FIRST}(\epsilon \vdash)] \end{array}$$

I'm using ϵ here to signify that β in the original item matched an empty string. $\text{FIRST}(\epsilon \vdash)$ is $\{\vdash\}$. Continue the closure operation by closing both of the newly added items: The first one is:

$$\begin{array}{c} [s \rightarrow \alpha \cdot x \beta, C] \\ \hline [e \rightarrow \cdot e + t, \vdash] \end{array}$$

so add:

$$\begin{array}{c} [s \rightarrow \cdot \gamma, \text{FIRST}(\beta C)] \\ [e \rightarrow \cdot e + t, \text{FIRST}(+t \vdash)] \\ [e \rightarrow \cdot t, \text{FIRST}(+t \vdash)] \end{array}$$

Note that these items differ from the previous ones only in the lookahead sets. Since $+$ is not nullable, $\text{FIRST}(+t \vdash)$ is $+$. The second production derived directly from the kernel item was:

$$\begin{array}{c} [s \rightarrow \alpha \cdot x \beta, C] \\ \hline [e \rightarrow \cdot t, \vdash] \end{array}$$

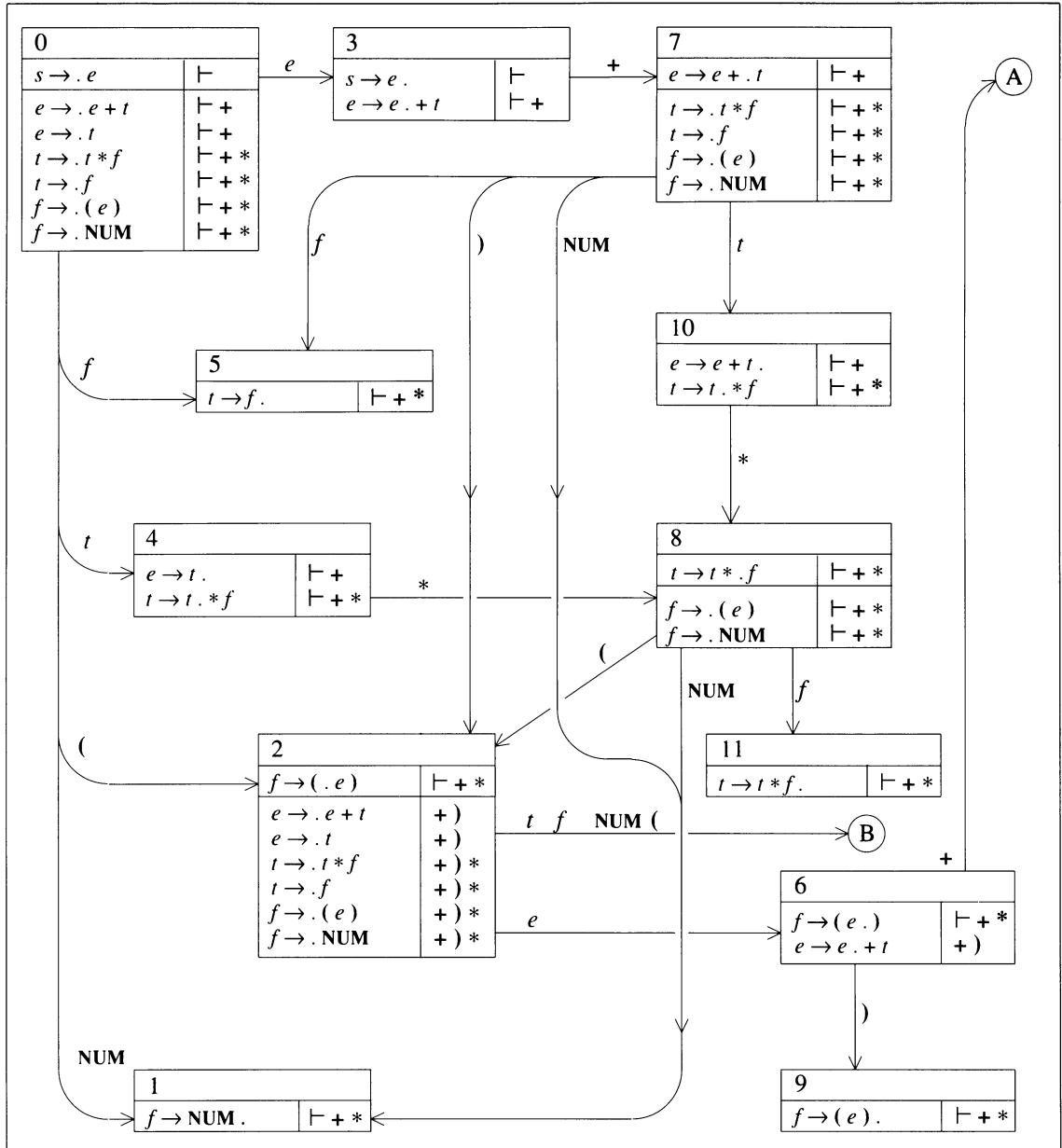
so add:

$$\begin{array}{c} [s \rightarrow \cdot \gamma, \text{FIRST}(\beta C)] \\ [t \rightarrow \cdot t * f, \text{FIRST}(\epsilon \vdash)] \\ [t \rightarrow \cdot f, \text{FIRST}(\epsilon \vdash)] \end{array}$$

The process continues in this manner until no more new LR(1) items can be created. The next states are created as before, adding edges for all symbols to the right of the dot and moving the dots in the kernel items of the new machine. The entire LR(1) state machine for our grammar is shown in Figure 5.7. I've saved space in the Figure by merging together all items in a state that differ only in lookaheads. The lookaheads for all such items are shown on a single line in the right column of each state. Figure 5.8 shows how the other closure items in State 0 are derived. Derivations for items in States 2 and 14 of the machine are also shown.

The LR(1) lookaheads are the relevant parts of the FOLLOW sets that we looked at earlier. A reduction from a given state is indicated only if the dot is at the far right of a production and the current lookahead symbol is in the LR(1) lookahead set for the item

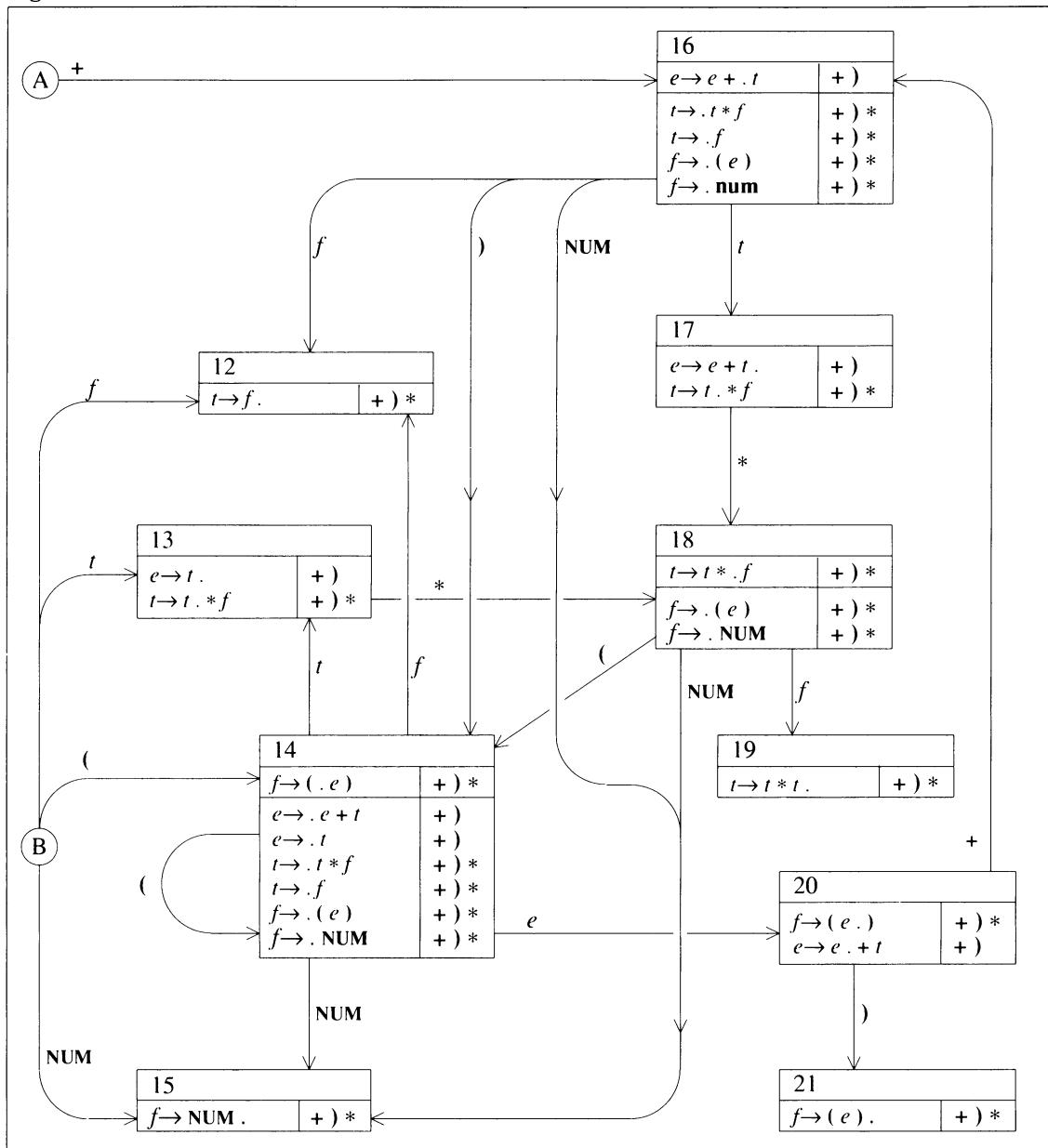
Figure 5.7. LR(1) State Machine



that contains that production. As before, a shift is indicated when the lookahead symbol matches a terminal label on an outgoing edge; outgoing edges marked with nonterminals are transitions in the goto part of the parse table. If there's a conflict between the lookahead sets and one of the outgoing edges, then the input grammar is not LR(1).

ϵ items.

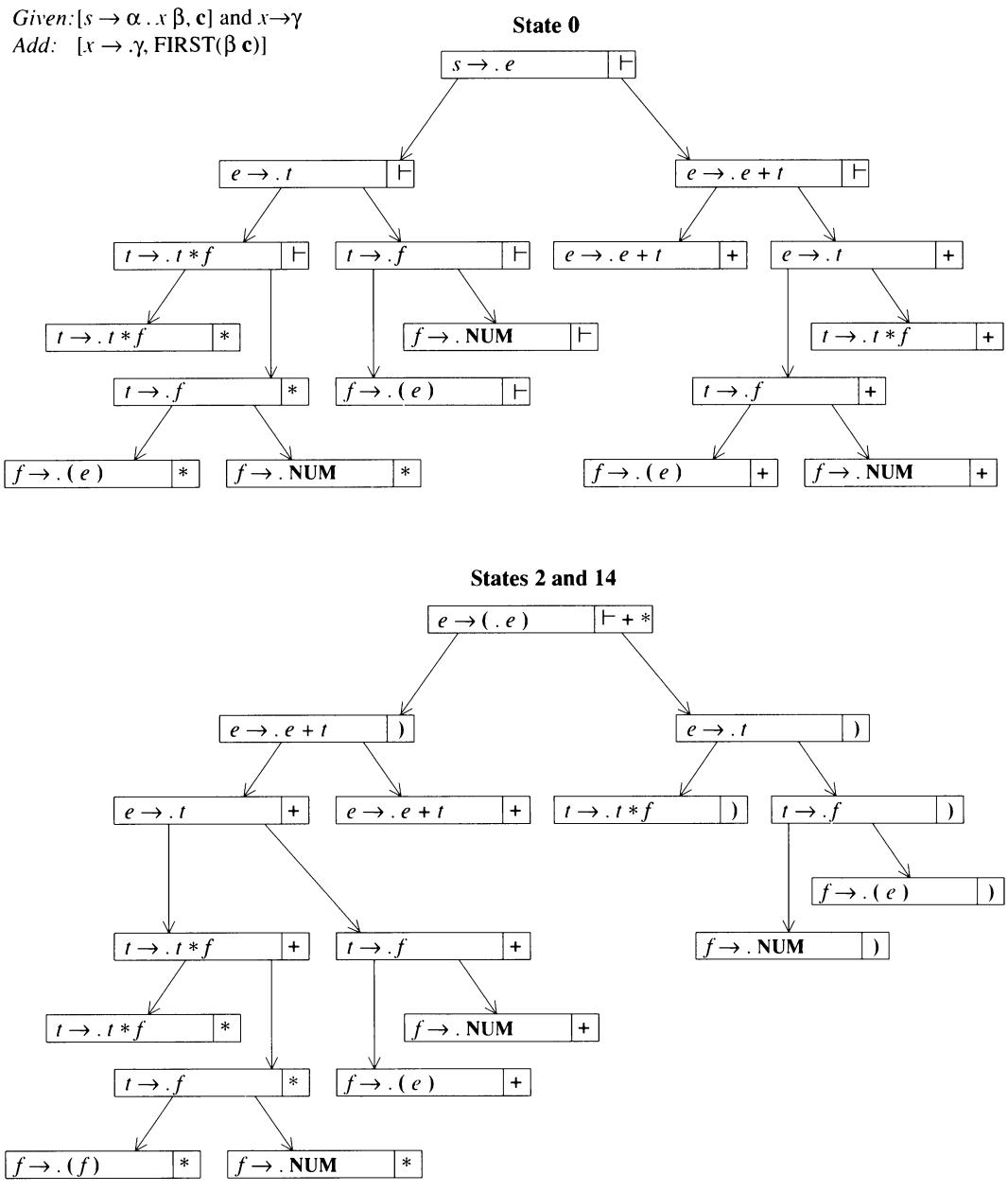
ϵ items are special only in that they're a degraded case of a normal LR(1) item. In particular, ϵ items take the form $[s \rightarrow \cdot \epsilon, C]$, where the lookahead C was inherited from the item's parent in the derivation tree. An ϵ item is never a kernel item, so will never generate transitions to a next state. The action associated with an ϵ item is always a reduce on C from the current state.

Figure 5.7. continued. LR(1) State Machine**5.6.4 LALR(1) Grammars***

The main disadvantage of a LR(1) state machine is that it is typically twice the size of an equivalent LR(0) machine, which is pretty big to start out with [the C grammar used in the next chapter uses 287 LR(0) states]. Fortunately, it's possible to have the best of both worlds. Examining the LR(1) machine in Figure 5.7, you are immediately struck by the similarity between the two halves of the machine. A glance at the lookahead tells you what's happening. A close parenthesis is a valid lookahead character only in a parenthesized subexpression—it can't appear in a lookahead set until after the open parenthesis has been processed. Similarly, the end-of-input marker is not a valid lookahead symbol inside a parenthesized expression because you must go past a close

Figure 5.8. Deriving Closure Items for LR(1) States 0, 2 and 14

Given: $[s \rightarrow \alpha . x \beta, c]$ and $x \rightarrow \gamma$
 Add: $[x \rightarrow .\gamma, \text{FIRST}(\beta c)]$

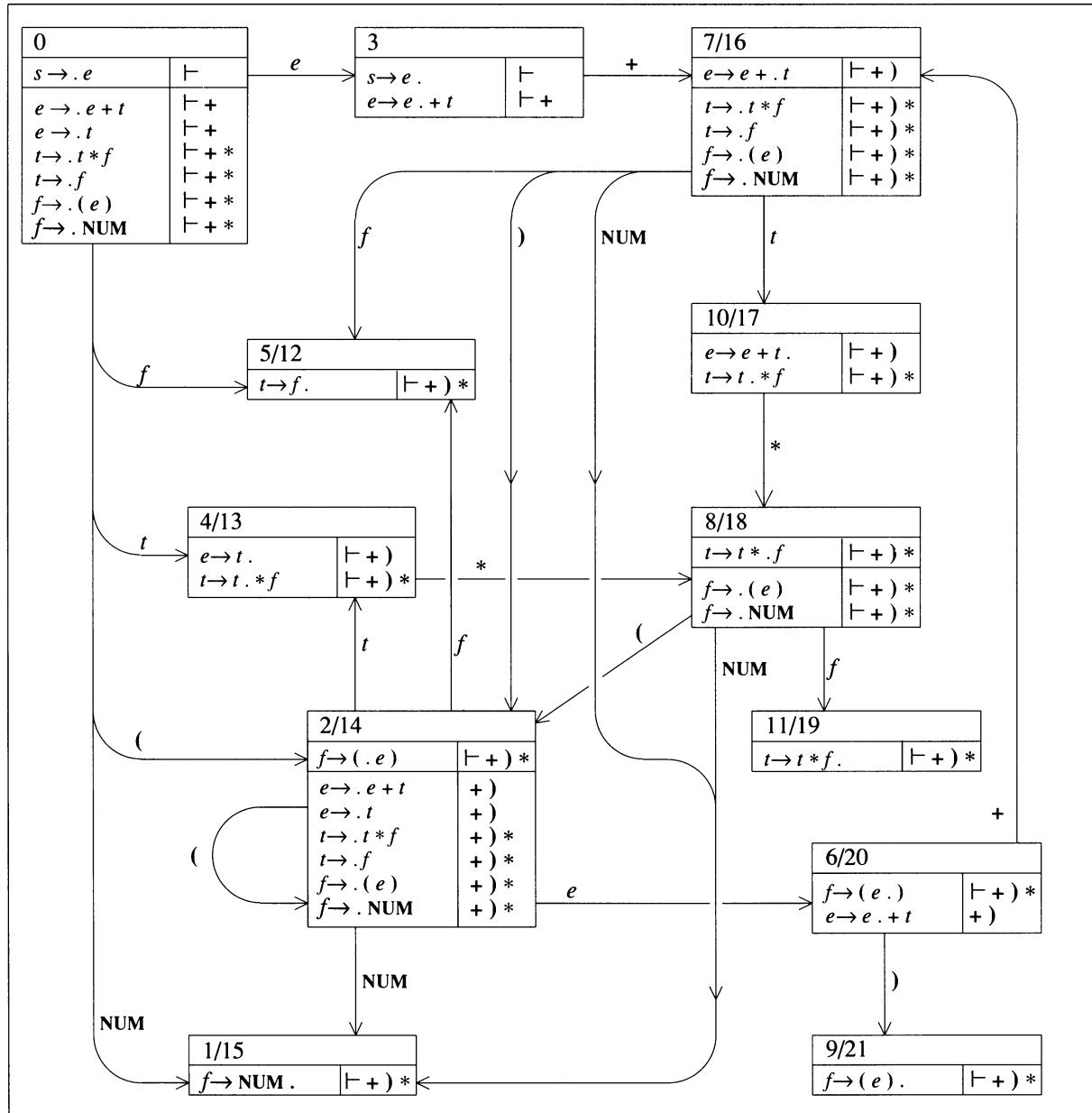


parenthesis first. The outer part of the machine (all of the left half of Figure 5.7 except States 6 and 9) handles unparenthesized expressions, and the inner part (States 6 and 9, and all of the right half of Figure 5.7) handles parenthesized subexpressions. The parser moves from one part to the other by reading an open parenthesis and entering State 2, which acts as a portal between the inner and outer machines. The parser moves back into the outer part of the machine by means of a reduction; it pops enough items so that the uncovered state number is in the outer machine.

You can take advantage of the similarity between the machines by merging them together. Two LR(1) states can be merged if the production and dot-position

components of all kernel items are the same. The states are merged by combining the lookahead sets for the items that have the same production and dot position, and then modifying the next-state transitions to point at the merged state. For example, States 2 and 14 are identical in all respects but the lookaheads. You can merge them together, forming a new State 2/14, and then go through the machine changing all next-state transitions to either State 2 or State 14 to point at the new State 2/14. This merging process creates the machine pictured in Figure 5.9.

Figure 5.9. LALR(1) State Machine for the Expression Grammar



The merging process does not affect error detection in the grammar. It usually takes a little longer for the merged machine to detect an error, but the merged machine reads no more input than the unmerged one. There is one pitfall, however. The merging does,

Merging does not affect error recovery, might add inadequate states.

after all, add elements to the lookahead sets. Consequently, there might be inadequate states in the merged machine that are not present in the original LR(1) machine. (This is rare, but it can happen.) If no such conflicts are created by the merging process, the input grammar is said to be an LALR(1) grammar.

Manufacturing the tables.

If you compare the LALR(1) machine in Figure 5.9 with the LR(0) machine in Figure 5.6 on page 360 you'll notice that the machines are identical in all respects but the lookahead characters. This similarity is no accident, and the similarity of the machines can be exploited when making the tables; you can save space when making an LALR(1) machine by constructing an LR(0) machine directly and adding lookaheads to it, rather than creating the larger LR(1) machine and then merging equivalent states. You can do this in several ways—the method used by *occs* is the simplest. It proceeds as if it were making the LR(1) machine, except that before creating a new state, it checks to see if an equivalent LR(0) state (an LR(1) state with the same productions and dot positions in the kernel) already exists. If it finds such a state, *occs* adds lookaheads to the existing state rather than creating a new one. I'll look at this procedure in depth, below.

More efficient table-generation methods.

There are other, more efficient methods for creating LALR(1) machines directly from the grammar that I won't discuss here. An alternate method that is essentially the same as the foregoing, but which uses different data structures and is somewhat faster as a consequence, is suggested in [Aho], pp. 240–244. An LR(0) state machine is created and then lookaheads are added to it directly. You can see the basis of the method by looking at Figure 5.8 on page 366. Lookaheads appear on the tree in one of two ways, either they propagate down the tree from a parent to a child or they appear spontaneously. A lookahead character propagates from the item $[s \rightarrow \alpha.x\beta, c]$ to its child when $\text{FIRST}(\beta)$ is nullable. It appears spontaneously when $\text{FIRST}(\beta)$ contains only terminal symbols. An even more efficient method is described in [DeRemer79] and [DeRemer82]; it is summarized in [Tremblay], pp. 375–383.

5.7 Representing LR State Tables

The LALR(1) state machine is easily represented by a two dimensional array of integers, where the value of the number represents the actions. The row index is typically the state number, and the columns are indexed by numbers representing the tokens and nonterminals. The LALR(1) state machine in Figure 5.9 is represented by the two-dimensional array in Table 5.11. (This is the same table as in Figure 5.9.) Shifts can be represented by positive numbers, and reductions can be represented with negative numbers. For example, a $s6$ directive can be represented by the number 6, and an $r5$ can be represented by a -5 . A very large integer that's not likely to be a state number can be used to signify error transitions. The accept action, which is really a reduction by Production 0, can be represented by a zero.

Positive numbers=shift,
Negative
numbers=reduce, 0=ac-
cept.

LALR(1) parse-table size.

Though the table can be represented by a two-dimensional array, such a table is likely to be very large. Taking the C grammar that's used in the next chapter as characteristic, the grammar has 182 productions, 50 terminal symbols, and 75 nonterminal symbols. It generates 310 LALR(1) states. The resulting LALR(1) transition matrix has as many rows as there are states, and as many columns as there are terminal and nonterminal symbols (plus one for the end-of-input marker), so there are $(50+75+1) = 125$ columns and 310 rows, or 38,750 cells. Given a 2-byte int, that's 77,500 bytes. In most systems, it's worthwhile to minimize the size of the table.

Pair-compression in
LALR(1) parse tables.

Only 3,837 cells (or 10%) of the 38,750 cells in the uncompressed table are nonerror transitions. As a consequence, you can get considerable compression by using the pair-compression method discussed in Chapter Two: Each row in the table is represented by an array, the first element of which is a count of the number of pairs in the array, and the

Table 5.11. Tabular Representation of Transition Diagram in Figure 5.9

	Shift/Reduce Table (Yy_action)					Goto Table (Yy_goto)			
	⊤	NUM	+	*	()	s	e	t	f
0	–	s1	–	–	s2	–	3	4	5
1	r6	–	r6	r6	–	r6	–	–	–
2	–	s1	–	–	s2	–	6	4	5
3	accept	–	s7	–	–	–	–	–	–
4	r2	–	r2	s8	–	r2	–	–	–
5	r4	–	r4	r4	–	r4	–	–	–
6	–	–	s7	–	–	s9	–	–	–
7	–	s1	–	–	s2	–	–	10	5
8	–	s1	–	–	s2	–	–	–	11
9	r5	–	r5	r5	–	r5	–	–	–
10	r1	–	r1	s8	–	r1	–	–	–
11	r3	–	r3	r3	–	r3	–	–	–
Error transitions are marked with –									

remainder of which is series of [nonterminal,action] pairs. For example, Row 4 of the transition matrix can be represented as follows:

```
state_4[] = { 4, [], r2], [+ , r2], [⊤, r2], [* , s8] };
```

The initial 4 says that there are four pairs that follow, the first pair says to reduce by Production 2 if the input symbol is a close parentheses, and so forth. If the current input symbol is not attached to any pair, then an error transition is indicated.

Another array (call it the *index* array) is indexed by state number and each cell evaluates to a pointer to one of the pairs arrays. If an element of the index array is NULL, then there are no outgoing transitions from the indicated state. If several rows of the table are identical (as is the case in the current table), only a single copy of the row needs to be kept, and several pointers in the index array can point at the single row. The *goto* portion of the table is organized the same way, except that the row array contains [nonterminal, next state] pairs. Figure 5.10 shows how the table in Listing 5.11 is represented in an *occs* output file. The actual code is in Listing 5.2.

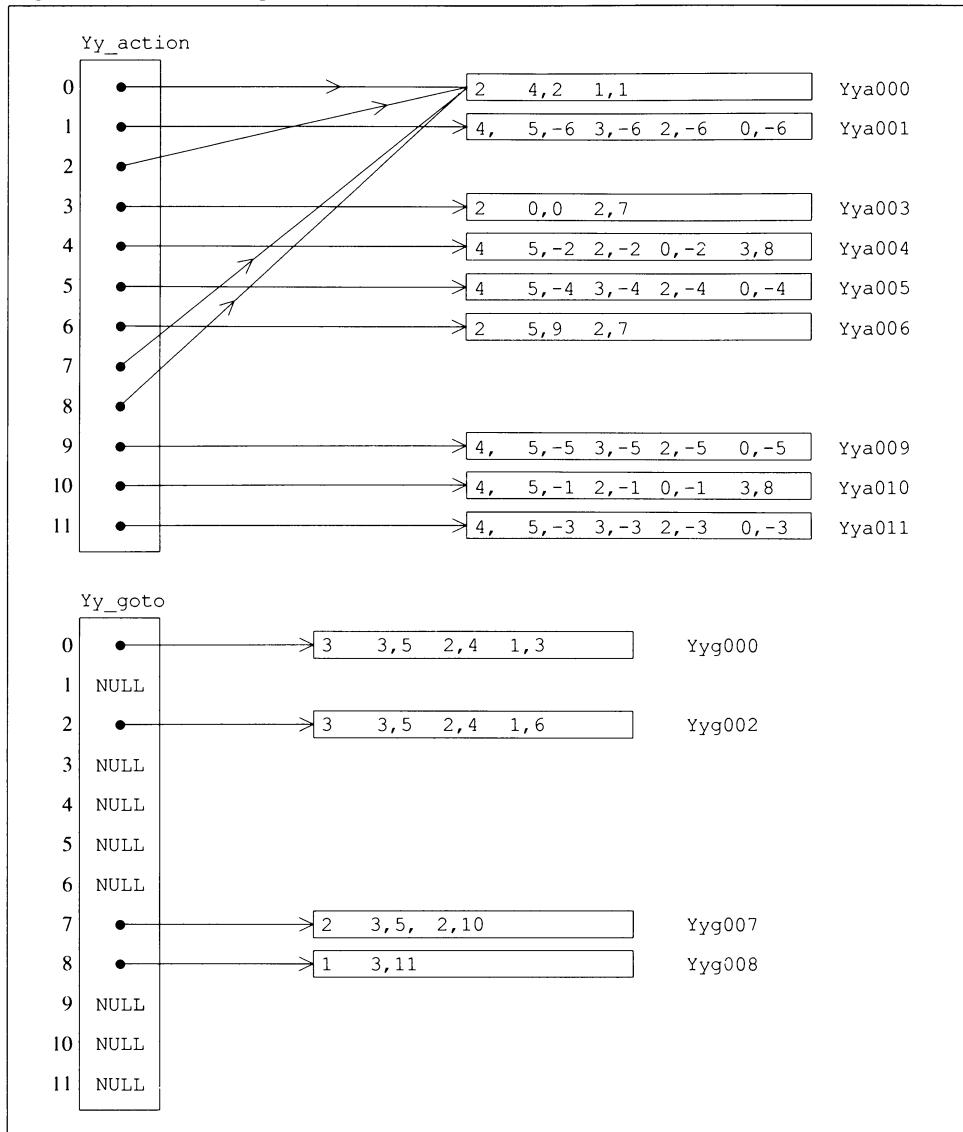
The *Yya* arrays on lines 20 to 28 of Listing 5.2 represent the rows, and *Yy_action* (on lines 30 to 34) is the index array; the pairs consist of tokenized input symbols and action descriptions, as described earlier. The values used for tokens are as follows:

YyaN arrays, Yy_action

Token	Value returned from lexical analyzer
⊤	0
NUM	1
+	2
*	3
(4
)	5

Looking at row four of the table (*Yya004* on line 23), the pair [5,–2] is a reduction by Production 2 (–2) on a close parenthesis (5). [3,8] is a shift to State 8 (8) on a * (3), and so on. The tables can be decompressed using the *yy_next()* subroutine in Listing 5.3.

Parse-table decompression.

Figure 5.10. Actual Representation of Parse Tables for Grammar in Listing 5.8**Listing 5.2.** C Representation of the Parse Tables

```

1  /*-----
2   * The Yy_action table is action part of the LALR(1) transition
3   * matrix. It's compressed and can be accessed using the yy_next()
4   * subroutine, declared below.
5   *
6   *           Yya000[]={ 3,      5, 3 ,     2, 2 ,     1, 1   };
7   * state number---+ | | |
8   * number of pairs in list---+ | |
9   * input symbol (terminal)-----+ |
10  * action-----+
11  *
12  * action = yy_next( Yy_action, cur_state, lookahead_symbol );
13  *
14  *       action < 0    -- Reduce by production n,  n == -action.

```



Listing 5.2. continued...

```

15     *      action == 0    -- Accept (ie. reduce by production 0)
16     *      action > 0    -- Shift to state n, n == action.
17     *      action == YYF -- error
18   */
19
20 int Yya000[]={ 2, 4,2 , 1,1 };
21 int Yya001[]={ 4, 5,-6 , 3,-6 , 2,-6 , 0,-6 };
22 int Yya003[]={ 2, 0,0 , 2,7 };
23 int Yya004[]={ 4, 5,-2 , 2,-2 , 0,-2 , 3,8 };
24 int Yya005[]={ 4, 5,-4 , 3,-4 , 2,-4 , 0,-4 };
25 int Yya006[]={ 2, 5,9 , 2,7 };
26 int Yya009[]={ 4, 5,-5 , 3,-5 , 2,-5 , 0,-5 };
27 int Yya010[]={ 4, 5,-1 , 2,-1 , 0,-1 , 3,8 };
28 int Yya011[]={ 4, 5,-3 , 3,-3 , 2,-3 , 0,-3 };
29
30 int *Yy_action[12] =
31 {
32     Yya000, Yya001, Yya000, Yya003, Yya004, Yya005, Yya006, Yya000,
33     Yya009, Yya010, Yya011
34 };
35
36 /*-----
37  * The Yy_goto table is goto part of the LALR(1) transition matrix. It's com-
38  * pressed and can be accessed using the yy_next() subroutine, declared below.
39  *
40  * nonterminal = Yy_lhs[ production number by which we just reduced ]
41  *
42  *           Yyg000[]={ 3, 5,3 , 2,2 , 1,1 };
43  * uncovered state--+
44  * number of pairs in list--+
45  * nonterminal-----+
46  * goto this state-----+
47  * goto_state = yy_next( Yy_goto, cur_state, nonterminal );
48 */
49
50 int Yyg000[]={ 3, 3,5 , 2,4 , 1,3 };
51 int Yyg002[]={ 3, 3,5 , 2,4 , 1,6 };
52 int Yyg007[]={ 2, 3,5 , 2,10 };
53 int Yyg008[]={ 1, 3,11 };
54
55 int *Yy_goto[12] =
56 {
57     Yyg000, NULL , Yyg002, NULL , NULL , NULL , NULL , Yyg007, Yyg008,
58     NULL , NULL , NULL
59 };
60
61 /*-----
62  * The Yy_lhs array is used for reductions. It is indexed by production
63  * number and holds the associated left-hand side adjusted so that the
64  * number can be used as an index into Yy_goto.
65  */
66
67 YYPRIVATE int Yy_lhs[7] =
68 {
69     /* 0 */      0,
70     /* 1 */      1,
71     /* 2 */      1,

```



Listing 5.2. continued...

```

72      /* 3 */      2,
73      /* 4 */      2,
74      /* 5 */      3,
75      /* 6 */      3
76  };
77
78 /*-----
79 * The Yy_reduce[] array is indexed by production number and holds
80 * the number of symbols on the right-hand side of the production.
81 */
82
83 YYPRIVATE int Yy_reduce[7] =
84 {
85     /* 0 */      1,
86     /* 1 */      3,
87     /* 2 */      1,
88     /* 3 */      3,
89     /* 4 */      1,
90     /* 5 */      3,
91     /* 6 */      1
92 }

```

Listing 5.3. yynext.c—Table Decompression

```

1 int yy_next( table, cur_state, symbol )
2 int **table;                                /* table to decompress */
3 int symbol;                                 /* column index */
4 int cur_state;                             /* row index */
5 {
6     /* Given current state and input symbol, return next state. */
7
8     int *p = table[ cur_state ] ;
9     int i;
10
11    if( p )
12        for( i = (int) *p++; --i >= 0 ; p += 2 )
13            if( symbol == p[0] )
14                return p[1];
15
16    return YYF;                               /* error indicator */
17 }

```

Goto transitions. `Yy_lhs`,
`Yy_reduce`.

Goto transitions could be
compressed by column.

The goto transitions are computed using two arrays. `Yy_lhs` (on line 67 of Listing 5.2) is indexed by production number and evaluates to an integer representing the symbol on the left-hand side of that production. That integer is, in turn, used to search the goto part of the array for the next-state transition. A second array, `Yy_reduce` on line 83 of Listing 5.2, is indexed by production number and evaluates to the number of symbols on the right-hand side of that production. A simplified version of the `occs` parser loop, which demonstrates how the tables are used, is shown in Listing 5.4.

Since there tend to be fewer nonterminals than states, the `Yy_goto` array could be made somewhat smaller if it were compressed by column rather than by row. The `Yy_goto` array could be indexed by nonterminal symbol and could contain `[current_state, next_state]` pairs. `OCCS` doesn't do this both because two table-decompression subroutines would be needed and because the parser would run more

Listing 5.4. Simplified Parser Loop Demonstrating Parse-Table Usage

```

1 push( 0 );                      /* Push start state */
2 lookahead = yylex();           /* get first lookahead symbol */
3
4 while( 1 )
5 {
6     do_this = yy_next( Yy_action, state_at_top_of_stack(), lookahead );
7
8     if( do_this == YYF )
9     {
10        /* ERROR */
11    }
12    else if( do_this > 0 )          /* Simple shift action */
13    {
14        lookahead = yylex();       /* advance */
15        push( do_this );         /* push new state */
16    }
17    else                          /* reduction */
18    {
19        yy_act( -do_this );      /* do code-generation action */
20
21        if( do_this == YYACCEPT )
22            break;
23        else
24        {
25            production_number = -do_this;
26
27            rhs_length = Yy_reduce[ production_number ];
28
29            while( --rhs_length >= 0 )          /* pop rhs_length items */
30                pop();
31
32            next_state = yy_next( Yy_goto, state_at_top_of_stack(),
33                                Yy_lhs[production_number] );
34            push( next_state );
35        }
36    }
37 }

```

slowly because the chain lengths would be a bit longer. You can also combine the two arrays, putting shift actions, reduce actions, and goto transitions together in a single row array; but again, it would take longer to decompress the table because of the extra chain lengths.

5.8 Eliminating Single-Reduction States*

Reexamining Table 5.11 on page 369, you'll notice that States 1, 5, 9, and 11 all consist of nothing but error transitions and reductions by a single production. For example, all the actions in State 1 are reductions by Production 6. These states, called *single-reduction states*, are typically caused by productions that are used only to add precedence rules to a grammar—they are created by productions with a single symbol on the right-hand side, such as the $e \rightarrow t$ and $t \rightarrow f$ in the grammar we've been using.

You can see what these states are actually doing by looking at how one of the states is used as the parse progresses. If a **NUM** is found in State 0, for example, the parser shifts to State 1, where the only legal action is a reduce by Production 6 if the next input

Single-reduction states.

character is \vdash , +, * or close parenthesis. This reduction just gets us back to State 0, without having advanced the input. All this backing and forthing is really unnecessary because the parser could reduce by Production 6 directly from State 0 if the next input character is a **NUM** and the following character is a \vdash , +, * or close parenthesis. If something else is in the input, the error is caught the next time the parser tries to shift, so the error information is actually duplicated in the table.

Removing single-reduction states from parse table.

All single-reduction states can be removed from the table by introducing another class of action items. If a row of the table contains a shift to a single-reduction state, replace it with a directive of the form dN , where N is the number of the production by which the parser would reduce had it gone to the single-reduction state. In the current table, for example, all $s1$ directives would be replaced by $d6$ directives. The parser processes a d directive just like a normal reduce directive, except that it pops one fewer than the number of symbols on the right-hand side of the indicated production, because it didn't push the last of these symbols.

A similar substitution has to be made in the *goto* portion of the table—all transitions into single-reduction states must be replaced with equivalent d directives. This change forces a second modification to the parser because several code-generation actions might now have to be executed during a reduction if a d directive is encountered in the *goto* table. The earlier table is modified to eliminate single-reduction states in Table 5.12. A parser algorithm for the new table is in Table 5.13.

Table 5.12. Parse Table with Single-Reduction States Removed

	Shift/Reduce Table (Yy_action)					Goto Table (Yy_goto)			
	\vdash	NUM	+	*	()	s	e	t	f
0	—	d6	—	—	s2	—	—	3	4 d4
2	—	d6	—	—	s2	—	—	6	4 d4
3	accept	—	s7	—	—	—	—	—	—
4	r2	—	r2	s8	—	r2	—	—	—
6	—	—	s7	—	—	d5	—	—	—
7	—	d6	—	—	s2	—	—	—	10 d4
8	—	d6	—	—	s2	—	—	—	d3
10	r1	—	r1	s8	—	r1	—	—	—
Error transitions are marked with —									

Can't remove single-reduction state if there is an action.

There's one final caveat. You cannot eliminate a single-reduction state if there is a code-generation action attached to the associated production because the stack will have one fewer items on it than it should when the action is performed—you won't be able to access the attributes correctly. In practice, this limitation is enough of a problem that **occs** doesn't use the technique. In any event, the disambiguating rules discussed in the next section eliminate many of the single-reduction states because the productions that cause them are no longer necessary.

There is one thing that's easy to do that does not affect the grammar at all, however, and can significantly reduce the table sizes in a pair-compressed array. Since the error information associated with single-reduction states is redundant (because the error is caught with the next shift), you can replace the entire chain of pairs that represent the single-reduction state with a single, default, reduce action that is performed regardless of the input symbol. The **UNIX yacc** utility takes this idea even further by providing a default action for every row of the table that represents the most common action in that row. The problem here is the error recovery, which is now very difficult to do because it is difficult to determine legal lookahead characters while doing the recovery.

Table 5.13. LR Parser Algorithm for Single-Reduction-Minimized Tables

```

push( 0 );
while( Action[TOS][input] ≠ Accept )
{
    if( Action[TOS][input] = - )
        error( );
    else if ( Action[TOS][input] = sX )
    {
        push( X );
        advance( );
    }
    else if ( Action[TOS][input] = rX )
    {
        act( X );
        pop( as many items as are in the RHS of production X );
        while( (i = Goto[ uncovered TOS ][ LHS of production X ]) = dX )
        {
            act( X );
            pop( one fewer than the number of symbols in the RHS of production X );
        }
        push i;
    }
    else if ( Action[TOS][input] = dX )
    {
        act( X );
        pop( one fewer than the number of symbols in the RHS of production X )
        while( (i = Goto[ uncovered TOS ][ LHS of production X ]) = dX )
        {
            act( X );
            pop( one fewer than the number of symbols in the RHS of production X );
        }
        push i;
    }
}
accept();

```

5.9 Using Ambiguous Grammars*

Even though LR(1) and LALR(1) grammars are more flexible than the SLR(1) and LR(0), they still can't be ambiguous, because, as you would expect, an ambiguous grammar yields an ambiguous parse. The ambiguity is reflected in the resulting state machine as a shift/reduce conflict. Ambiguous productions are, nonetheless, useful. They tend to make a grammar both smaller and easier to read. The two most common ambiguities are found in Productions 2, 4, and 5 of the following grammar:

1.	<i>stmt</i>	\rightarrow	<i>expr SEMI</i>
2.			IF <i>stmt ELSE stmt</i>
3.			IF <i>stmt</i>
4.	<i>expr</i>	\rightarrow	<i>expr PLUS expr</i>
5.			<i>expr TIMES expr</i>
6.			ID
ID is an identifier			

Resolving ambiguities
when making parse
tables.

This grammar creates the state machine pictured in Figure 5.11, and the ambiguities are reflected as shift/reduce conflicts in the States 5, 10, and 11. The LALR(1) lookahead heads are shown in brackets next to those items that trigger reductions.

Though the foregoing grammar is not LALR(1), it is nonetheless possible to use it to create a parser. The basic strategy examines the various shift/reduce conflicts in the machine, and makes decisions as to whether a shift or reduce is appropriate in that situation, basing the decision on the semantic rather than syntactic description of the language. In practical terms, you need to know the relative precedence and the associativity of those tokens that appear in ambiguous productions, and this information can be used to disambiguate the tables.

Precedence and associativity information is used to resolve ambiguities in arithmetic expressions as follows: State 10 contains the following items:

$$\begin{array}{lcl} \textit{expr} & \rightarrow & \textit{expr} \cdot + \textit{expr} \\ \textit{expr} & \rightarrow & \textit{expr} + \textit{expr} \cdot \quad [; + *] \\ \textit{expr} & \rightarrow & \textit{expr} \cdot * \textit{expr} \end{array}$$

The parser wants to reduce by

$$\textit{expr} \rightarrow \textit{expr} + \textit{expr}$$

if the lookahead symbol is a semicolon, plus sign, or star; it also wants to shift on a plus sign or star. The semicolon poses no problem, but the arithmetic operators do.

Now, consider the parse caused by the input *a+b+c*. The stack and input looks like this after the *a*, *+*, and *b* are processed:

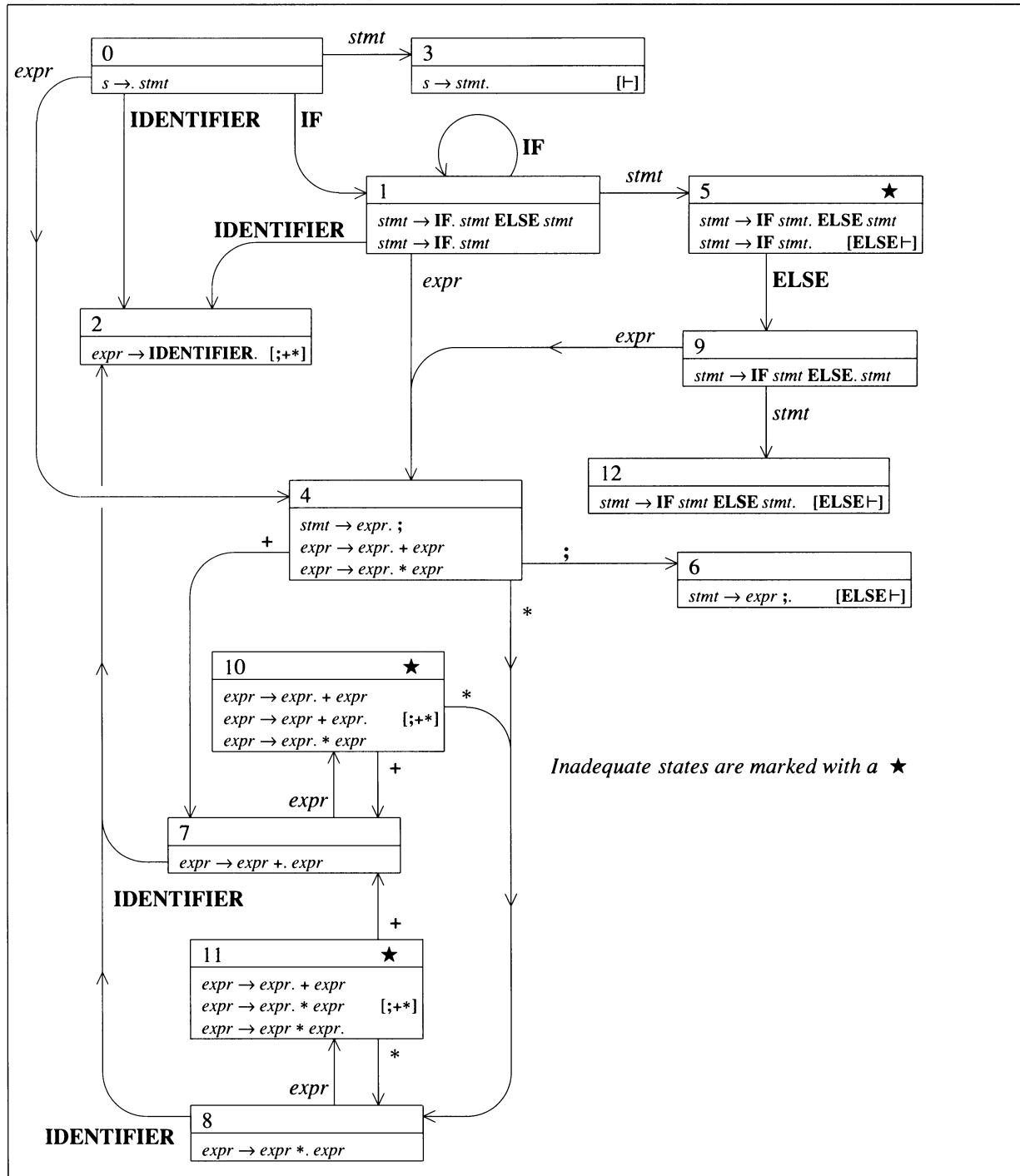
<i>expr + expr</i>	+ c ;
--------------------	-------

The parser now looks ahead and sees the plus, and it has a problem. If the plus operator is left associative, it will want to reduce immediately (before processing the next plus sign). A left-associative parse would continue as follows:

<i>expr + expr</i>	+ c ;
<i>expr</i>	+ c ;
<i>expr +</i>	c ;
<i>expr + ID</i>	;
<i>expr + expr</i>	;
<i>expr</i>	;
<i>expr ;</i>	
<i>stmt</i>	

This way the left subexpression would be processed before going on with the input. Right associativity is handled differently. Here the parser wants to put the whole expression onto the stack so that it can reduce from right to left, like this:

Figure 5.11. State Machine for an Ambiguous Grammar



<i>expr + expr</i>	+ c ;
<i>expr + expr +</i>	c ;
<i>expr + expr + ID</i>	;
<i>expr + expr + expr</i>	;
<i>expr + expr</i>	;
<i>expr</i>	;
<i>expr ;</i>	
<i>stmt</i>	

So the decision of whether to shift or reduce here is based on the desired associativity of the plus operator. In this case, you get left associativity by resolving the conflict in favor of the reduce.

Shift for higher precedence, reduce for lower.

State 10 has one other problem—illustrated by the input $a+b*c$. Here, the star is higher precedence than the plus, so the multiplication must be done first. The conflict is active when the following stack condition exists:

<i>expr + expr</i>	* c ;
--------------------	-------

Reducing at this point would effectively ignore the relative precedence of + and *; the addition would be done first. Resolving in favor of the shift, however, correctly generates the following parse:

<i>expr + expr *</i>	c ;
<i>expr + expr * ID</i>	;
<i>expr + expr * expr</i>	;
<i>expr + expr</i>	;
<i>expr</i>	;
<i>expr ;</i>	
<i>stmt</i>	

and the multiplication is done before the addition. So, the conflict should be resolved in favor of the shift if the lookahead operator is higher in precedence than the token closest to the top of the stack.

Dangling else.

Precedence and associativity help only if a production implements an arithmetic expression, however. The sensible default for handling a shift/reduce conflict in nonarithmetic productions is to decide in favor of the shift. This decision takes care of the conflict in State 5, which contains the following items

IF stmt . ELSE stmt
IF stmt . [← ELSE]

Here, the machine wants both to shift and reduce when the lookahead token is an **ELSE**. The problem is exemplified by the following input:

```

IF
    ID ;
ELSE
    IF
        ID ;
ELSE
    ID ;

```

Most programming languages require an **ELSE** to bind with the closest preceding **IF**. A partial parse of the foregoing generates the following stack:

IF stmt ELSE IF ID ; ELSE ID ;

A shift is required in this situation, because an **ELSE** is present in the input. Were you to reduce **IF stmt** to **stmt**, the **ELSE** would be left dangling in the input. That is, since a **stmt** can't start with an **ELSE**, a syntax error would be created if the parser reduced.

Letting the parse continue—shifting on every conflict—you'll eventually get the following on the stack:

IF stmt ELSE IF stmt ELSE stmt

which can then be reduced to:

IF stmt ELSE stmt

and then to:

stmt

Since the rightmost statement is reduced first, the **ELSE** binds to the closest **IF**, as required. The foregoing rules can be summarized as follows:

- (1) Assign precedence and associativity information to all terminal symbols and precedence information to productions.⁴ If no precedence and associativity is assigned, a terminal symbol is assigned a precedence level of zero (very low) and is nonassociative. Productions are assigned the same precedence level as the rightmost terminal symbol in the production or zero if there are no terminal symbols on the right-hand side.⁵
- (2) When a LALR(1) shift/reduce conflict is encountered, the precedence of the terminal symbol to be shifted is compared with the precedence of the production by which you want to reduce. If either the terminal or the production is of precedence zero, then resolve the conflict in favor of the shift. Otherwise, if the precedences are equal, resolve using the following table:

Rules for resolving ambiguities in parse table.

associativity of lookahead symbol	resolve in favor of
left	reduce
right	shift
nonassociative	shift

Otherwise, if the precedences are not equal, use the following table:

4. **Yacc** and **occs** do this with the `%left`, `%right` and `%nonassoc` directives.

5. **Yacc** and **occs** let you override this default with a `%prec TOKEN` directive where `TOKEN` is a terminal symbol that was declared with a previous `%left`, `%right`, or `%nonassoc`. The production is assigned the same precedence level as the token in the `%prec` directive. Note that the `%prec` directive must be to the left of the semicolon or vertical bar that terminates the production.

precedence	resolve in favor of
lookahead symbol < production	reduce
lookahead symbol > production	shift

Resolving reduce/reduce conflicts.

There's another kind of conflict that can appear in your grammar—a *reduce/reduce* conflict. Consider the following grammar that implements one rule of Kernighan and Cherry's troff preprocessor, **eqn**.⁶

1.	<i>expr</i>	\rightarrow	<i>expr SUB expr SUP expr</i>
2.			<i>expr SUB expr</i>
3.			<i>expr SUP expr</i>
4.			<i>LB expr RB</i>
5.			CONST

Eqn typesets equations. The implemented grammar does the following:

Input	Typeset as follows:
X sub 1	X ₁
X sup 2	X ²
X sub 1 sup 2	X ₁ ²

The last case has to be handled specially to prevent the subscript and superscript from getting out of line, like this:

X₁²

The earlier grammar takes care of the special case by adding an extra production (Production 1) that takes care of a case that would actually be handled by the rest of the grammar were it not there—X sub 1 sup 2 would be parsed without difficulty if Production 1 wasn't in the grammar. It wouldn't be recognized as a special case, however. The grammar as presented yields a state with the following kernel items:

<i>expr</i>	\rightarrow	<i>expr SUB expr SUP expr .</i>	[\vdash SUB SUP RB]
<i>expr</i>	\rightarrow	<i>expr SUP expr .</i>	[\vdash SUB SUP RB]
<i>expr</i>	\rightarrow	<i>expr . SUB expr SUP expr</i>	
<i>expr</i>	\rightarrow	<i>expr . SUB expr</i>	
<i>expr</i>	\rightarrow	<i>expr . SUP expr</i>	

Here, the parser can't decide which of the two productions that have dots on the far right should be reduced. (This situation is a *reduce/reduce* conflict.) The sensible default is to resolve the conflict in favor of the production that is higher in the grammar.

Note that *reduce/reduce* conflicts should generally be avoided by modifying the grammar. More often than not, this type of conflict is caused by an error in the grammar, like this:

<i>foo</i>	\rightarrow	TOKEN
	...	
		TOKEN

6. This example is taken from [Aho], pp. 252–254. It's significant that neither I nor any of the other compiler programmers that I know could think of a real example of a valid *reduce/reduce* conflict other than the current one. *Reduce/reduce* conflicts are almost always caused by errors in the grammar.

They can also appear when you do seemingly reasonable things like this, however:

<i>karamazov</i>	→	<i>dmitry</i> <i>ivan</i> <i>alyosha</i>
<i>dmitry</i>	→	NAME <i>do_something</i>
<i>ivan</i>	→	NAME <i>do_something_else</i>
<i>alyosha</i>	→	NAME <i>do_yet_another_thing</i>

This second situation is essentially the same as the earlier example. Several of *karamazov*'s right-hand sides begin with the same terminal symbol. Here, however, the fact has been disguised by an intervening nonterminal. Because of the first-production-has-higher-precedence rule, *ivan* and *alyosha* are never executed.⁷ The solution is to restructure the grammar as follows:

<i>karamazov</i>	→	NAME <i>brother</i>
<i>brother</i>	→	<i>dmitry</i> <i>ivan</i> <i>alyosha</i>
<i>dmitry</i>	→	<i>do_something</i>
<i>ivan</i>	→	<i>do_something_else</i>
<i>alyosha</i>	→	<i>do_yet_another_thing</i>

or to go a step further:

<i>karamazov</i>	→	NAME <i>brother</i>
<i>brother</i>	→	<i>do_something</i>
		<i>do_something_else</i>
		<i>do_yet_another_thing</i>

5.10 Implementing an LALR(1) Parser—The `occ`s Output File

The remainder of this chapter discusses both the `occ`s output file and how `occ`s itself works; you can skip to the next chapter if you're not interested in this level of detail. In any event, you should read Appendix E, which contains a user's manual for `occ`s, before continuing, because `occ`s will be used in the next chapter.

Most of the code that comprises `occ`s is taken directly from **LLama**, and was discussed in the last chapter. The two main differences are the parser itself and the parse-table generation routines. This section discusses the parser, using the output file created from the `occ`s input file in Listing 5.5. (The associated lexical analyzer is defined by the `LEX` input file in Listing 5.6.)

The grammar used here is not a particularly good use of `occ`s, because I've built precedence and associativity information into the grammar itself, rather than using the `%left` and `%right` directives as I did with the expression grammar in Appendix E. I've done this because `occ`s fudges the tables for ambiguous grammars, using the rules described in the previous section. For the purposes of discussion, however, it's better to work with a grammar that is a strict LALR(1) grammar. Note that I've extended the definition of the `NUM` token to recognize an identifier as well as a number (and renamed it to `NUM_OR_ID` as a consequence).

7. In `yacc` and `occ`s, be careful of blithely using `%prec` to solve a problem like the foregoing without considering the consequences. The problem may well have been caused by a deficiency in the grammar and, though the `%prec` may eliminate the warning message, the problem will still exist.

Listing 5.5. *expr.y*—Odds Input File for an Expression Compiler

```

1 %term NUM_OR_ID      /* a number or identifier */ */
2
3 %left PLUS           /* +          (lowest precedence) */ */
4 %left STAR            /* *          */ */
5 %left LP RP           /* ( )        (highest precedence) */ */
6
7 %}
8 #include <stdio.h>
9 #include <ctype.h>
10 #include <malloc.h>
11 #include <tools/debug.h>
12 #include <tools/stack.h>
13
14 extern char *yytext;
15
16 stack_dcl( Namepool, char*, 10 );           /* Stack of 10 temporary-var names */
17 #define freename(x) push( Namepool, (x) ) /* Release a temporary variable */
18 #define getname()    pop ( Namepool )     /* Allocate an temporary variable */
19
20 typedef char *stype;           /* Value stack is stack of char pointers */
21 #define YYSTYPE stype
22
23 #define YYSHIFTACT(tos) (*tos="")           /* Shift a null string */
24 %}
25
26 %%
27 /* A small expression grammar that recognizes numbers, names, addition (+),
28 * multiplication (*), and parentheses. Expressions associate left to right
29 * unless parentheses force it to go otherwise. * is higher precedence than +.
30 */
31
32 s      : e ;
33
34 e      : e PLUS t { yycode("%s += %s\n", $1, $3); freename($3); }
35 | t    /* $$ = $1 */
36 ;
37
38 t      : t STAR f { yycode("%s *= %s\n", $1, $3); freename($3); }
39 | f    /* $$ = $1 */
40 ;
41
42 f      : LP e RP { $$ = $2; }
43 | NUM_OR_ID {
44             /* Copy operand to a temporary. Note that I'm adding an
45             * underscore to external names so that they can't con-
46             * flict with the compiler-generated temporary names
47             * (t0, t1, etc.). */
48
49             yycode("%s = %s%s\n", $$ = getname(),
50                   isdigit(*yytext) ? "" : "_",
51                   yytext );
52         }
53 ;
54
55 %%
56 /*-----*/
57

```

Listing 5.5. continued...

```

58 char      *yypstk( vptr, dptr )
59 char      **vptr;                      /* Value-stack pointer */
60 char      *dptr;                      /* Symbol-stack pointer (not used) */
61 {
62     /* Yystk is used by the debugging routines. It is passed a pointer to a
63     * value-stack item and should return a string representing that item. Since
64     * the current value stack is a stack of string pointers, all it has to do
65     * is dereference one level of indirection.
66     */
67
68     return *vptr ? *vptr : "-";
69 }
70
71 /*-----*/
72
73 yy_init_occ()
74 {
75     /* Called by yyparse just before it starts parsing. Initialize the
76     * temporary-variable-name stack and output declarations for the variables.
77     */
78
79     push( Namepool, "t9" ); push( Namepool, "t8" ); push( Namepool, "t7" );
80     push( Namepool, "t6" ); push( Namepool, "t5" ); push( Namepool, "t4" );
81     push( Namepool, "t3" ); push( Namepool, "t2" ); push( Namepool, "t1" );
82     push( Namepool, "t0" );
83
84     yycode( "public word t0, t1, t2, t3, t4;\n" );
85     yycode( "public word t5, t6, t7, t8, t9;\n" );
86 }
87
88 /*-----*/
89
90 main( argc, argv )
91 char      **argv;
92 {
93     yy_get_args( argc, argv );
94
95     if( argc < 2 )
96         ferr("Need file name\n");
97
98     else if( ii_newfile(argv[1]) < 0 )
99         ferr( "Can't open %s\n", argv[1] );
100
101    yyparse();
102    exit( 0 );
103 }

```

Occs generates several output files for these input files. A symbol-table dump is in Listing 5.7 (*yyout.sym*), and token definitions are in *yyout.h*, Listing 5.8. The symbol table shows the internal values used for other various symbols including the production numbers that are assigned by **occ**s to the individual productions. These numbers are useful for deciphering the tables and debugging. The token-definitions file is typically **#included** in the L^EX input file so that the analyzer knows which values to return. The token names are taken directly from the **%term** and **%left** declarations in the **occ**s input file.

Occs token definitions
(*yyout.h*), symbol table
(*yyout.sym*).

Listing 5.6. *expr.lex*—LeX Input File for an Expression Compiler

```

1  %
2  #include "yyout.h"
3  %
4  digit    [0-9]
5  alpha     [a-zA-Z_]
6  alnum    [0-9a-zA-Z_]
7  %%
8
9  "+"          return PLUS;
10 " *"         return STAR;
11 "("          return LP;
12 ")"          return RP;
13 {digit}+     |
14 {alpha}{alnum}* return NUM_OR_ID;
15 .
16 %%

```

Listing 5.7. *yyout.sym*—Symbol Table Generated from Input File in Listing 5.5

```

1 ----- Symbol table -----
2
3 NONTERMINAL SYMBOLS:
4 e (257)  <>
5   FIRST : NUM_OR_ID LP
6     1: e -> e PLUS t .....PREC 1
7     2: e -> t
8
9 f (259)  <>
10  FIRST : NUM_OR_ID LP
11    5: f -> LP e RP .....PREC 3
12    6: f -> NUM_OR_ID
13
14 s (256)  (goal symbol) <>
15  FIRST : NUM_OR_ID LP
16    0: s -> e
17
18 t (258)  <>
19  FIRST : NUM_OR_ID LP
20    3: t -> t STAR f .....PREC 2
21    4: t -> f
22
23
24 TERMINAL SYMBOLS:
25
26 name      value  prec  assoc  field
27 LP        4      3      1      <>
28 NUM_OR_ID 1      0      -      <>
29 PLUS     2      1      1      <>
30 RP       5      3      1      <>
31 STAR     3      2      1      <>

```

Listing 5.8. *yyout.h*— Token Definitions Generated from Input File in Listing 5.5

```

1 #define _EOI_      0
2 #define NUM_OR_ID  1
3 #define PLUS       2
4 #define STAR       3
5 #define LP         4
6 #define RP         5

```

The *occ*s-generated parser starts in Listing 5.9. As before, listings for those parts of the output file that are just copied from the template file are labeled *occ.par*. The *occ*s-generated parts of the output file (tables, and so forth) are in listings labeled *yyout.c*. Since we are really looking at a single output file, however, line numbers carry from one listing to the other, regardless of the name. The output starts with the file header in Listing 5.9. Global variables likely to be used by yourself (such as the output streams) are defined here. Note that *<stdio.h>*, *<stdarg.h>*, and *<tools/yystack.h>*, are **#included** on lines one to three. *<stdarg.h>* contains macros that implement the ANSI variable-argument mechanism and *<tools/yystack.h>* contains stack-maintenance macros. These last two files are both described in Appendix A. The remainder of the header, in Listing 5.10, is copied from the header portion of the input file.

Listing 5.9. *occ.par*— File Header

```

1 #include <stdio.h>
2 #include <stdarg.h>
3 #include <tools/yystack.h>
4
5 FILE *yycodeout = stdout ;           /* Output stream (code). */
6 FILE *yybssout = stdout ;           /* Output stream (bss). */
7 FILE *yydataout = stdout ;          /* Output stream (data). */
8 int yylookahead ;                  /* Lookahead token. */
9
10 extern char *yytext;                /* Declared by lex in lexyy.c */
11 extern int yylineno;
12 extern int yyleng;
13
14 extern char *ii_ptext() ;           /* Lookback function used by lex */
15 extern int ii_plength() ;          /* in /src/compiler/lib/input.c. */
16 extern int ii_plineno() ;
17
18 #ifdef YYDEBUG                   /* Define YYD here so that it can be used */
19 # define YYD(x) x                /* in the user-supplied header. */
20 #else
21 # define YYD(x) /* empty */
22 #endif
23
24 /*-----*/

```

The output file continues in Listing 5.11 with various macro definitions. As was the case in the LEX output file, these macros are not redefined if you defined them in the header part of the input file. (That's what all the **#ifndefs** are for.) Note that `printf()` calls are mapped to `yycode()` calls on line 44 if debugging is enabled. I've done this because the standard `printf()` circumvents the window-output mechanism and messes up the screen. This way, all `printf()` calls in the *occ*s input file, at least, will behave properly. Direct output to the screen in other files is still a problem,

Macro definitions.

`printf()` mapped to
`yycode()`.

Listing 5.10. *yyout.c*—Code Taken from Input-Specification Header

```

25 #include <stdio.h>
26 #include <ctype.h>
27 #include <malloc.h>
28 #include <tools/debug.h>
29 #include <tools/stack.h>
30
31 extern char *yytext;
32
33 stk_dcl( Namepool, char*, 10 );           /* Stack of 10 temporary-var names */
34 #define freename(x) push( Namepool, (x) ) /* Release a temporary variable */
35 #define getname()    pop ( Namepool )      /* Allocate a temporary variable */
36
37 typedef char          *stype;           /* Value stack is stack of char pointers */
38 #define YYSTYPE        stype
39
40 #define YYSHIFTACT(tos) (*tos="")           /* Shift a null string */

```

Stack-macro customization.

however. The stack macros, discussed in Appendix A and **#included** on line three of Listing 5.9, are customized on line 87 of Listing 5.11. The redefinition of *yystk_cls* causes all the stacks declared with subsequent *yystk_dcl()* invocations to be **static**.

Listing 5.11. *occsp.par*—Definitions

```

41 #undef YYD                      /* Redefine YYD in case YYDEBUG was defined */
42 #ifdef YYDEBUG                  /* explicitly in the header rather than with */
43 # define YYD(x) x              /* a -D on the occs command line. */
44 # define printf yycode         /* Make printf() calls go to output window */
45 #else
46 # define YYD(x) /* empty */
47 #endif
48
49 #ifndef YYACCEPT
50 # define YYACCEPT return(0)   /* Action taken when input is accepted. */
51 #endif
52
53 #ifndef YYABORT
54 # define YYABORT return(1)   /* Action taken when input is rejected. */
55 #endif
56
57 #ifndef YYPRIVATE
58 # define YYPRIVATE static    /* define to a null string to make public */
59 #endif
60
61 #ifndef YYMAXERR
62 # define YYMAXERR 25        /* Abort after this many errors */
63 #endif
64
65 #ifndef YYMAXDEPTH
66 # define YYMAXDEPTH 128     /* State and value stack depth */
67 #endif
68
69 #ifndef YYCASCADE
70 # define YYCASCADE 5       /* Suppress error msgs. for this many cycles */
71 #endif
72

```



Listing 5.11. continued...

```

73 #ifndef YYSTYPE          /* Default value stack type */          */
74 # define YYSTYPE int
75 #endif
76                                     /* Default shift action: inherit $$ */ */
77 #ifndef YYSHIFTACT
78 # define YYSHIFTACT(tos) ( (tos)[0] = yylval )
79 #endif
80
81 #ifdef YYVERBOSE
82 # define YYV(x) x
83 #else
84 # define YYV(x)
85 #endif
86
87 #undef yystk_cls           /* redefine stack macros for local */
88 #define yystk_cls YYPRIVATE      /* use. */                      */
89
90 /*
91 * #defines used in the tables. Note that the parsing algorithm assumes that
92 * the start state is State 0. Consequently, since the start state is shifted
93 * only once when we start up the parser, we can use 0 to signify an accept.
94 * This is handy in practice because an accept is, by definition, a reduction
95 * into the start state. Consequently, a YYR(0) in the parse table represents an
96 * accepting action and the table-generation code doesn't have to treat the
97 * accepting action any differently than a normal reduce.
98 *
99 * Note that if you change YY_TTYPE to something other than short, you can no
100 * longer use the -T command-line switch.
101 */
102
103 #define YY_IS_ACCEPT    0          /* Accepting action (reduce by 0) */
104 #define YY_IS_SHIFT(s) ((s) > 0)   /* s is a shift action */        */
105
106 typedef short    YY_TTYPE;
107 #define YYF      ((YY_TTYPE)( (unsigned short)^0 >>1 ))
108
109 /*
110 * Various global variables used by the parser. They're here because they can
111 * be referenced by the user-supplied actions, which follow these definitions.
112 *
113 * If -p or -a was given to OCCS, make Yy_rhslen and Yy_val (the right-hand
114 * side length and the value used for $$) public, regardless of the value of
115 * YYPRIVATE (yylval is always public). Note that occs generates extern
116 * statements for these in yyacts.c (following the definitions section).
117 */
118
119 #if !defined(YYACTION) || !defined(YYPARSER)
120 # define YYP /* nothing */
121 #else
122 # define YYP    YYPRIVATE
123 #endif
124
125 YYPRIVATE int     yynerrs = 0;          /* Number of errors. */        */
126
127 yystk_dcl( Yy_stack, int, YYMAXDEPTH ); /* State stack. */            */
128
129     YYSTYPE yylval;                  /* Attribute for last token. */
130 YYP YYSTYPE Yy_val;                /* Used to hold $$.

```



Listing 5.11. continued...

```

131 YYP YYSTYPE Yy_vstack[ YYMAXDEPTH ];
132 YYP YYSTYPE *Yy_vsp;
133
134 YYP int Yy_rhslen;
135
136
/* Value stack. Can't use      */
/* yystack.h macros because   */
/* YYSTYPE could be a struct.* */
/* Number of nonterminals on  */
/* right-hand side of the     */
/* production being reduced. */

```

YY_IS_SHIFT,
YY_TTYPE.

The definitions on lines 103 to 107 of Listing 5.11 are used for the tables. Shift actions are represented as positive numbers (the number is the next state), reduce operations are negative numbers (the absolute value of the number is the production by which you're reducing) and zero represents an accept action (the input is complete when you reduce by Production 0). **YY_IS_SHIFT** on line 104 is used to differentiate between these. **YY_TTYPE** on the next line is the table type. You should probably change it to **short** if your machine uses a 16-bit **short** and 32-bit **int**. **YY_TTYPE** must be signed, and a **char** is usually too small because of the number of states in the machine.

Error marker: **YYF**.

YYF, on line 107 of Listing 5.11, represents failure transitions in the parse tables. (**YYF** is not stored in the compressed table, but is returned by the table-decompression subroutine, **yy_act_next()**, which I'll discuss in a moment.) It evaluates to the largest positive **short int** (with two's complement numbers). Breaking the macro down: **(unsigned)** **~0** is an unsigned int with all its bits set. The **unsigned** suppresses sign extension on the right shift of one bit, which yields a number with all but the high bit set. The resulting quantity is cast back to a **YY_TTYPE** so that it will agree with other elements of the table. Note that this macro is not particularly portable, and might have to be changed if you change the **YY_TTYPE** definition on the previous line. Just be sure that **YYF** has a value that can't be confused with a normal shift or reduce directive.

YYP.

The final part of Listing 5.11 comprises declarations for parser-related variables that might be accessed in one of the actions. The state and value stacks are defined here (along with the stack pointers), as well as some house-keeping variables. Note that those variables of class **YYP** are made public if **-a** or **-p** is specified to **occs**. (In which case, **YYACTION** and **YPARSER** are not both present—the definition is output by **occs** itself at the top of the file, and the test is on line 119).

Action subroutine,
yy_act().

Listing 5.12 holds the action subroutine, which executes the code-generation actions from the **occs** input file. Various tables that **occs** generates from the input grammar are also in this listing. The actions imbedded in the input grammar are output as **case** statements in the **switch** in **yy_act()** (on lines 137 to 169 of Listing 5.12). As in **LLama**, the **case** values are the production numbers. Each production is assigned a unique but arbitrary number by **occs**; these numbers can be found in **yyout.sym** in Listing 5.7 on page 384, which is generated when **occs** finds a **-D**, **-s**, or **-S** command-line switch. The production numbers precede each production in the symbol-table output file. The start production is always Production 0, the next one in the input file is Production 1, and so forth.

Translated dollar attributes: **\$\$**, **\$1**, etc.

Note that the dollar attributes have all been translated to references to the value stack at this juncture. For example, on line 164, the line:

```
t : t STAR f { yycode("%s *= %s\n", $1, $3); freename($3); }
```

has generated:

```
{ yycode("%s *= %s\n", yyvsp[2], yyvsp[0]); freename(yyvsp[0]); }
```

in the action subroutine. **yyvsp** is the value-stack pointer and a downward-growing stack is used. (A push is a ***--yyvsp=x**; a pop is a ***yyvsp++**.) The situation is

Listing 5.12. *yyout.c*—The Action Subroutine and Tables

```

137 yy_act( yy_production_number, yyvsp )
138     int     yy_production_number;
139     YYSTYPE *yyvsp;
140 {
141     /* This subroutine holds all the actions in the original input
142      * specification. It normally returns 0, but if any of your actions return a
143      * non-zero number, then the parser will halt immediately, returning that
144      * nonzero number to the calling subroutine. I've violated my usual naming
145      * conventions about local variables so that this routine can be put into a
146      * separate file by occs.
147     */
148
149     switch( yy_production_number )
150     {
151     case 1:
152         { yycode("%s += %s\n", yyvsp[2], yyvsp[0]); freename(yyvsp[0]); }
153         break;
154     case 6:
155         { yycode("%s = %s%s\n", Yy_val = getname(),
156                 isdigit(*yytext) ? "" : "_",
157                 yytext );
158         }
159         break;
160     case 5:
161         { Yy_val = yyvsp[1]; }
162         break;
163     case 3:
164         { yycode("%s *= %s\n", yyvsp[2], yyvsp[0]); freename(yyvsp[0]); }
165         break;
166     default: break; /* In case there are no actions */
167     }
168     return 0;
169 }
170
171 /-----
172     * Yy_stok[] is used for debugging and error messages. It is indexed
173     * by the internal value used for a token (as used for a column index in
174     * the transition matrix) and evaluates to a string naming that token.
175     */
176
177 char *Yy_stok[] =
178 {
179     /* 0 */    "_EOI_",
180     /* 1 */    "NUM_OR_ID",
181     /* 2 */    "PLUS",
182     /* 3 */    "STAR",
183     /* 4 */    "LP",
184     /* 5 */    "RP"
185 };
186
187 /-----
188     * The Yy_action table is action part of the LALR(1) transition matrix. It's
189     * compressed and can be accessed using the yy_next() subroutine, below.
190     *
191     *          Yya000[]={ 3, 5, 3 , 2, 2 , 1, 1 };
192     * state number---+ | | |
193     * number of pairs in list---+ | |
194     * input symbol (terminal)-----+ |
195     * action-----+

```



Listing 5.12. continued...

```

196  *
197  *     action = yy_next( Yy_action, cur_state, lookahead_symbol );
198  *
199  *     action < 0 -- Reduce by production n, n == -action.
200  *     action == 0 -- Accept (ie. reduce by production 0)
201  *     action > 0 -- Shift to state n, n == action.
202  *     action == YYF -- error
203  */
204
205 YYPRIVATE YY_TTYPE Yya000[]={ 2, 4, 2 , 1, 1 };
206 YYPRIVATE YY_TTYPE Yya001[]={ 4, 5,-6 , 3,-6 , 2,-6 , 0,-6 };
207 YYPRIVATE YY_TTYPE Yya003[]={ 2, 0, 0 , 2, 7 };
208 YYPRIVATE YY_TTYPE Yya004[]={ 4, 5,-2 , 2,-2 , 0,-2 , 3, 8 };
209 YYPRIVATE YY_TTYPE Yya005[]={ 4, 5,-4 , 3,-4 , 2,-4 , 0,-4 };
210 YYPRIVATE YY_TTYPE Yya006[]={ 2, 5, 9 , 2, 7 };
211 YYPRIVATE YY_TTYPE Yya009[]={ 4, 5,-5 , 3,-5 , 2,-5 , 0,-5 };
212 YYPRIVATE YY_TTYPE Yya010[]={ 4, 5,-1 , 2,-1 , 0,-1 , 3, 8 };
213 YYPRIVATE YY_TTYPE Yya011[]={ 4, 5,-3 , 3,-3 , 2,-3 , 0,-3 };
214
215 YYPRIVATE YY_TTYPE *Yy_action[12] =
216 {
217     Yya000, Yya001, Yya000, Yya003, Yya004, Yya005, Yya006, Yya000, Yya000,
218     Yya009, Yya010, Yya011
219 };
220
221 /-----
222 * The Yy_goto table is goto part of the LALR(1) transition matrix. It's com-
223 * pressed and can be accessed using the yy_next() subroutine, declared below.
224 *
225 * nonterminal = Yy_lhs[ production number by which we just reduced ]
226 *
227 *         Yyg000[]={ 3, 5, 3 , 2, 2 , 1, 1 };
228 * uncovered state--+
229 * number of pairs in list--+
230 * nonterminal-----+
231 * goto this state-----+
232 * goto_state = yy_next( Yy_goto, cur_state, nonterminal );
233 */
234
235 YYPRIVATE YY_TTYPE Yyg000[]={ 3, 3, 5 , 2, 4 , 1, 3 };
236 YYPRIVATE YY_TTYPE Yyg002[]={ 3, 3, 5 , 2, 4 , 1, 6 };
237 YYPRIVATE YY_TTYPE Yyg007[]={ 2, 3, 5 , 2, 10 };
238 YYPRIVATE YY_TTYPE Yyg008[]={ 1, 3, 11 };
239
240 YYPRIVATE YY_TTYPE *Yy_goto[12] =
241 {
242     Yyg000, NULL , Yyg002, NULL , NULL , NULL , Yyg007, Yyg008,
243     NULL , NULL , NULL
244 };
245
246 /-----
247 * The Yy_lhs array is used for reductions. It is indexed by production number
248 * and holds the associated left-hand side, adjusted so that the number can be
249 * used as an index into Yy_goto.
250 */
251
252 YYPRIVATE int Yy_lhs[7] =
253 {

```



Listing 5.12. continued...

```

254      /* 0 */      0,
255      /* 1 */      1,
256      /* 2 */      1,
257      /* 3 */      2,
258      /* 4 */      2,
259      /* 5 */      3,
260      /* 6 */      3
261 } ;
262
263 /-----
264 * The Yy_reduce[] array is indexed by production number and holds
265 * the number of symbols on the right hand side of the production.
266 */
267
268 YYPRIVATE int Yy_reduce[7] =
269 {
270     /* 0 */      1,
271     /* 1 */      3,
272     /* 2 */      1,
273     /* 3 */      3,
274     /* 4 */      1,
275     /* 5 */      3,
276     /* 6 */      1
277 };
278 #ifdef YYDEBUG
279
280 /-----
281 * Yy_slhs[] is a debugging version of Yy_lhs[]. Indexed by production number,
282 * it evaluates to a string representing the left-hand side of the production.
283 */
284
285 YYPRIVATE char *Yy_slhs[7] =
286 {
287     /* 0 */      "s",
288     /* 1 */      "e",
289     /* 2 */      "e",
290     /* 3 */      "t",
291     /* 4 */      "t",
292     /* 5 */      "f",
293     /* 6 */      "f"
294 };
295
296 /-----
297 * Yy_srhs[] is also used for debugging. It is indexed by production number
298 * and evaluates to a string representing the right-hand side of the production.
299 */
300
301 YYPRIVATE char *Yy_srhs[7] =
302 {
303     /* 0 */      "e",
304     /* 1 */      "e PLUS t",
305     /* 2 */      "t",
306     /* 3 */      "t STAR f",
307     /* 4 */      "f",
308     /* 5 */      "LP e RP",
309     /* 6 */      "NUM_OR_ID"
310 };
311 #endif

```

complicated by the fact that attributes are numbered from left to right, but the rightmost (not the leftmost) symbol is at the top of the parse stack. Consequently, the number that is part of the dollar attribute can't be used directly as an offset from the top of stack. You can use the size of the right-hand side to compute the correct offset, however. When a reduction is triggered, the symbols on the parse stack exactly match those on the right-hand side of the production. Given a production like $t \rightarrow t \text{ STAR } f$, f is at top of stack, **STAR** is just under the f , and t is under that. The situation is illustrated in Figure 5.12. yyvsp , the stack pointer, points at f , so $\$3$ translates to $\text{yyvsp}[0]$ in this case. Similarly, $\$2$ translates to $\text{yyvsp}[1]$, and $\$1$ translates to $\text{yyvsp}[2]$. The stack offset for an attribute is the number of symbols on the right-hand side of the current production less the number that is part of the dollar attribute. $\$1$ would be at $\text{yyvsp}[3]$ if the right-hand side had four symbols in it.

Figure 5.12. Translating $\$N$ to Stack References

$t \rightarrow t \text{ STAR } f$		<i>in rules section</i>	<i>in section following second %%</i>
Yy_vsp	→	$\$3 \quad \text{Yy_vsp}[0]$	$\text{Yy_vsp}[\text{Yy_rhslen}-3]$
		$\$2 \quad \text{Yy_vsp}[1]$	$\text{Yy_vsp}[\text{Yy_rhslen}-2]$
		$\$1 \quad \text{Yy_vsp}[2]$	$\text{Yy_vsp}[\text{Yy_rhslen}-1]$

Attributes in code section
of input file:

Right-hand-side length:
 Yy_rhslen .

Token-to-string conversion:
 $\text{Yy_stok}[]$.

$\text{Yy_slhs}[], \text{Yy_srhs}[]$.

Figure 5.12 also shows how attributes are handled when they are found in the third part of the **occs** input file rather than imbedded in a production. The problem here is that the number of symbols on the right-hand side is available only when **occs** knows which production is being reduced (as is the case when an action in a production is being processed). Code in the third section of the input file is isolated from the actual production, so **occs** can't determine which production the attribute references. The parser solves the problem by setting a global variable, Yy_rhslen , to the number of symbols on the right-hand side of the production being reduced. Yy_rhslen is modified just before each reduction. This variable can then be used at run-time to get the correct value-stack item. Note that negative attribute numbers are also handled correctly by **occs**.⁸ Figure 5.13 shows the value stack just before a reduction by $b \rightarrow d \ e \ f$ in the following grammar:

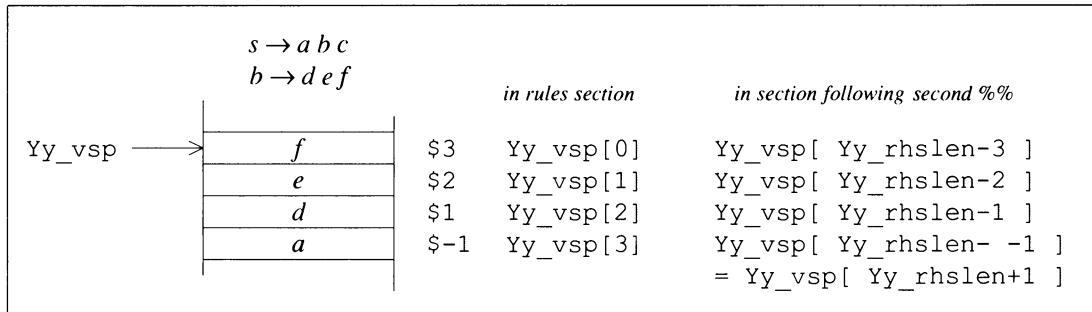
$$\begin{array}{ll} s & \rightarrow \ a \ b \ c \\ b & \rightarrow \ d \ e \ f \ \{x=\$-1;\} \end{array}$$

The $\$-1$ in the second production references the a in the partially assembled first production.

Listing 5.12 continues on line 177 with a token-to-string translation table. It is indexed by token value (as found in *yyout.h*) and evaluates to a string naming that token. It's useful both for debugging and for printing error messages. The other conversion tables are used only for the debugging environment, so are **#ifdefed** out when **YYDEBUG** is not defined. $\text{Yy_slhs}[]$ on line 285 is indexed by production number and holds a string representing the left-hand side of that production. $\text{Yy_srhs}[]$ on line 301 is

8. They aren't accepted by **yacc**.

Figure 5.13. Translating $\$-N$ to Stack References



similar, but it holds strings representing the right-hand sides of the productions.

Lines 187 to 277 of Listing 5.12 hold the actual parse tables. The state machine for our current grammar was presented earlier (in Table 5.11 on page 369) and the compressed tables were discussed in that section as well. Single-reduction states have not been eliminated here. I'll demonstrate how the parse tables are used with an example parse of $(1+2)$. The parser starts up by pushing the number of the start state, State 0, onto the state stack. It determines the next action by using the state number at the top of stack and the current input (lookahead) symbol. The input symbol is a left parenthesis (an LP token), which is defined in *yyout.h* to have a numeric value of 4. (It's also listed in *yyout.sym*.) The parser, then, looks in *Yy_action[0][4]*. Row 0 of the table is represented in *Yya000* (on Line 205 of Listing 5.12), and it's interested in the first pair: [4,2]. The 4 is the column number, and the 2 is the parser directive. Since 2 is positive, this is a shift action: 2 is pushed onto the parse stack and the input is advanced. Note that *Yya000* also represents rows 2, 6, and 7 of the table because all four rows have the same contents.

Using occs' compressed
parse tables.

The next token is a number (a `NUM_OR_ID` token, defined as 1 in `yyout.h`). The parser now looks at `Yy_action[2][1]` (2 is the current state, 1 the input token). Row 2 is also represented by `Yya000` on line 205 of Listing 5.12, and Column 1 is represented by the second pair in the list: [1,1]. The action here is a shift to State 1, so a 1 is pushed and the input is advanced again.

0 (1+2)

A 1 is now on the top of the stack, and the input symbol is a **PLUS**, which has the value 2. `Yy_action[1][2]` holds a -6 (it's in the third pair in `Yya000`, on line 205), which is a reduce-by-Production-6 directive (reduce by $f \rightarrow \text{NUM_OR_ID}$). The first thing the parser does is perform the associated action, with a `yy_act(6)` call. The actual reduction is done next. `Yy_reduce[6]` evaluates to the number of objects on the right-hand side of Production 6—in this case, 1. So, one object is popped, uncovering the previous-state number (2). Next, the goto component of the reduce operation is performed. The parser does this with two table lookups. First, it finds out which left-hand side is associated with Production 6 by looking it up in `Yy_lhs[6]` (it finds a 3 there). It then looks up the next state in `Yy_goto`. The machine is in State 2, so the row array is fetched from `Yy_goto[2]`, which holds a pointer to `Yyg002`. The parser then searches the pairs in `Yyg002` for the left-hand side that it just got from `Yy_lhs` (the 3), and it finds the pair [3,5]—the next state is State 5 (a 5 is pushed). The parse continues in this manner until a reduction by Production 0 occurs.

02 1+2)

Listing 5.13 contains the table-decompression routine mentioned earlier (`yy_next()` on line one), and various output subroutines as well. Two versions of the output routines are presented at the end of the listing, one for debugging and another version for production mode. The output routines are mapped to window-output functions if debugging mode is enabled, as in **LLama**. Also of interest is a third, symbol stack

0 2 1 +2)

[02] +2)

025 +2)

Table decompression:
yy_next().

Symbol stack:
Yy_dstack.

(Yy_dstack) defined on line 24. This stack is the one that's displayed in the middle of the debugging environment's *stack* window. It holds strings representing the symbols that are on the parse stack.

Listing 5.13. *occs.par*—Table-Decompression and Output Subroutines

```

1  YYPRIVATE YY_TTYPE yy_next( table, cur_state, inp )
2  YY_TTYPE **table;
3  YY_TTYPE cur_state;
4  int     inp;
5  {
6      /* Next-state routine for the compressed tables. Given current state and
7      * input symbol (inp), return next state.
8      */
9
10     YY_TTYPE *p = table[ cur_state ] ;
11     int     i;
12
13     if( p )
14         for( i = (int) *p++; --i >= 0 ; p += 2 )
15             if( inp == p[0] )
16                 return p[1];
17
18     return YYF;
19 }
20
21 /*-----*/
22 #ifdef YYDEBUG
23
24 yystk_dcl( Yy_dstack, char*, YYMAXDEPTH );           /* Symbol stack */
25
26 yycode( fmt )
27 char   *fmt;
28 {
29     va_list args;
30     va_start( args, fmt );
31     yy_output( 0, fmt, args );
32 }
33
34 yydata( fmt )
35 char   *fmt;
36 {
37     va_list args;
38     va_start( args, fmt );
39     yy_output( 1, fmt, args );
40 }
41
42 yybss( fmt )
43 char   *fmt;
44 {
45     va_list args;
46     va_start( args, fmt );
47     yy_output( 2, fmt, args );
48 }
49
50 /* yycomment() and yyerror() are defined in yydebug.c */
51
52 #else /*-----*/
53

```



Listing 5.13. continued...

```
54 #    define yy_nextoken()      yylex()      /* when YYDEBUG isn't defined. */
55 #    define yy_quit_debug()
56 #    define yy_init_debug()
57 #    define yy_pstack(x,y)
58 #    define yy_sym()
59
60 /* Use the following routines just like printf() to create output. The only
61 * differences are that yycode is sent to the stream called yycodeout, yydata
62 * goes to yydataout, and yybss goes to yybssout. All of these are initialized
63 * to stdout. It's up to you to close the streams after the parser terminates.
64 */
65
66 yycode( fmt )
67 char *fmt;
68 {
69     va_list args;
70     va_start( args,      fmt );
71     vfprintf( yycodeout, fmt, args );
72 }
73
74 yydata( fmt )
75 char *fmt;
76 {
77     va_list      args;
78     va_start( args,  fmt );
79     vfprintf( yydataout, fmt, args );
80 }
81
82 yybss( fmt )
83 char *fmt;
84 {
85     va_list      args;
86     va_start( args,  fmt );
87     vfprintf( yybssout, fmt, args );
88 }
89
90 void yycomment( fmt )
91 char *fmt;
92 {
93     va_list args;
94     va_start( args,  fmt );
95     vfprintf( stdout,  fmt, args );
96 }
97
98 void yyerror( fmt, ... )
99 char *
100 {
101     extern char *yytext;
102     extern int yylineno;
103     va_list args;
104
105     va_start( args, fmt );
106     fprintf ( stderr, "ERROR (line %d near %s): ", yylineno, yytext );
107     vfprintf( stderr, fmt, args );
108     fprintf ( stderr, "\n" );
109 }
110 #endif
```

Shift subroutine,
yy_shift().

The parser itself starts in Listing 5.14. The shift and reduce operations have been broken out into subroutines (on lines 111 and 128) in order to simplify the code in the parser itself. A new state is pushed onto the state stack on line 115. Garbage is pushed onto the value stack on line 116, but the user-supplied, default shift action [in the macro YYSHIFTACT()] is performed on the next line. This macro is passed a pointer to the top of the value stack (after the push)—the default action (defined on line 78 of Listing 5.11, page 387) pushes the contents of yyval onto the stack, so you can push an attribute for a nonterminal by modifying yyval from an accepting action in a LEX input file. (This process is discussed in Appendix E.) The debugging stack is kept aligned with the other stacks by pushing a string representing the shifted token on line 121. The yy_pstack() call on the next line updates the stack window and waits for another command if necessary. Breakpoints are also activated there.

Listing 5.14. *occs.par*—The Parser

```

111 YYPRIVATE    yy_shift( new_state, lookahead )      /* shift          */
112 int           new_state;                          /* push this state */
113 int           lookahead;                         /* Current lookahead */
114 {
115     yypush( Yy_stack, new_state );
116     --Yy_vsp;                                /* Push garbage onto value stack */
117     YYSHIFTACT( Yy_vsp );                     /* Then do default action */
118
119 #ifdef YYDEBUG
120     yycomment( "Shift %0.16s (%d)\n", Yy_stok[ lookahead ], new_state );
121     yypush_( Yy_dstack, Yy_stok[lookahead] );
122     yy_pstack(0, 1);
123 #endif
124 }
125
126 /*-----*/
127
128 YYPRIVATE      YY_reduce( prod_num, amount )
129 int            prod_num;                      /* Reduce by this production */
130 int            amount;                       /* # symbols on right-hand side of prod_num */
131 {
132     int next_state;
133
134     yypopn( Yy_stack, amount );                /* Pop n items off the state stack */
135     Yy_vsp += amount;                         /* and the value stack. */
136     *--Yy_vsp = Yy_val;                      /* Push $$ onto value stack */
137
138     next_state = yy_next( Yy_goto, yystk_item(Yy_stack,0), Yy_lhs[prod_num] );
139
140 #ifndef YYDEBUG
141
142     yypush_ ( Yy_stack, next_state );
143
144 #else
145
146     yy_break( prod_num );                    /* activate production breakpoint */
147
148     yypopn_ ( Yy_dstack, amount );
149
150     YYV( yycomment("    pop %d item%s\n", amount, amount==1 ? "" : "s"); )
151     yy_pstack( 0, 0 );
152

```

Listing 5.14. continued...

```

153     yypush_ ( Yy_stack, next_state );
154     yypush_ ( Yy_dstack, Yy_slhs[ prod_num ] );
155
156     YYV( yycomment("    push %0.16s (%d)", Yy_slhs[prod_num], next_state ); )
157
158     yy_pstack ( 0, 1 );
159
160 #endif
161 }
162
163 /*-----*/
164
165 YYPRIVATE void yy_init_stack() /* Initialize the stacks */
166 {
167     yystk_clear( Yy_stack );
168     yypush_ ( Yy_stack, 0 ); /* State stack = 0 */
169
170     Yy_vsp = Yy_vstack + (YYMAXDEPTH-1); /* Value stack = garbage */
171
172 # ifdef YYDEBUG
173     yystk_clear ( Yy_dstack );
174     yypush_ ( Yy_dstack, "$" );
175     yycomment ( "Shift start state\n" );
176     yy_pstack (0, 1); /* refresh stack window */
177 # endif
178 }
179
180 /*-----*/
181
182 YYPRIVATE int yy_recover( tok, suppress )
183
184 int tok; /* token that caused the error */
185 int suppress; /* No error message is printed if true */
186 {
187     int *old_sp = yystk_p(Yy_stack); /* State-stack pointer */
188     YYD( char **old_DSP = yystk_p(Yy_dstack); )
189     YYD( char *tos; )
190
191     if( !suppress )
192     {
193         yyerror( "Unexpected %s\n", Yy_stok[tok] );
194         if( ++yynerrs > YYMAXERR )
195         {
196             yyerror("Too many errors, aborting\n");
197             return 0;
198         }
199     }
200
201     do {
202
203         while( !yystk_empty(Yy_stack)
204             && yy_next( Yy_action, yystk_item(Yy_stack,0), tok ) == YYF )
205         {
206             yypop_( Yy_stack );
207
208             YYD( tos = yypop_(Yy_dstack); )
209             YYD( yycomment("Popping %d from state stack\n", tos); )
210             YYD( yy_pstack(0, 1); )
211         }
212     }

```



Listing 5.14. continued...

```

213     if( !yystk_empty(Yy_stack) )           /* Recovered successfully */
214     {
215         /* Align the value (and debug) stack to agree with the current
216         * state-stack pointer.
217         */
218
219         Yy_vsp = Yy_vstack + (yystk_ele(Yy_stack) - 1);
220
221 #ifdef YYDEBUG
222     yystk_p(Yy_dstack) = Yy_dstack + (yystk_ele(Yy_stack) - 1);
223     yycomment("Error recovery successful\n");
224     yy_pstack(0, 1);
225 #endif
226
227     return tok;
228 }
229
230     yystk_p( Yy_stack ) = old_sp ;
231
232     YYD( yystk_p( Yy_dstack ) = old_dsp ; )           )
233     YYD( yycomment("Restoring state stack."); )        )
234     YYD( yy_pstack(1, 1); )                           )
235     YYD( yycomment("discarding %s\n", Yy_stok[tok]); )  )
236
237 } while( ii_mark_prev(), tok = yy_nextoken() );
238
239 YYD( yycomment("Error recovery failed\n"); )
240 return 0;
241 }
242
243 /*-----*/
244
245 int      yyparse()
246 {
247     /* General-purpose LALR parser. Return 0 normally or -1 if the error
248     * recovery fails. Any other value is supplied by the user as a return
249     * statement in an action.
250     */
251
252     int act_num;          /* Contents of current parse table entry      */
253     int errcode;          /* Error code returned from yy_act()          */
254     int tchar;            /* Used to \0-terminate the lexeme           */
255     int suppress_err;    /* Set to YYCASCADE after an error is found   */
256                             /* and decremented on each parse cycle. Error */
257                             /* messages aren't printed if it's true.      */
258
259 #ifdef YYDEBUG
260     if( !yy_init_debug( Yy_stack, &yystk_p(Yy_stack),
261                         Yy_dstack, &yystk_p(Yy_dstack),
262                         Yy_vstack, sizeof(YYSTYPE), YYMAXDEPTH) )
263         YYABORT;
264 #endif
265
266     yy_init_stack();          /* Initialize parse stack                      */
267     yy_init_occs( Yy_vsp );
268
269     yylookahead = yy_nextoken(); /* Get first input symbol                   */
270     suppress_err = 0;
271

```



Listing 5.14. continued...

```

272     while( 1 )
273     {
274         act_num = yy_next( Yy_action, yystk_item(Yy_stack,0), yylookahead );
275
276         if( suppress_err )
277             --suppress_err;
278
279         if( act_num == YYF )
280         {
281             if( ! (yylookahead = yy_recover( yylookahead, suppress_err )) )
282                 YYABORT;
283
284             suppress_err = YYCASCADE;
285         }
286         else if( YY_IS_SHIFT(act_num) )           /* Simple shift action */
287         {
288             /* Note that yytext and yyleng are undefined at this point because
289              * they were modified in the else clause, below. You must use
290              * ii_text(), etc., to put them into a reasonable condition if
291              * you expect to access them in a YY_SHIFT action.
292             */
293
294             yy_shift( act_num, yylookahead );
295
296             ii_mark_prev();
297             yylookahead = yy_nextoken();
298         }
299         else
300         {
301             /* Do a reduction by -act_num. The activity at 1, below, gives YACC
302              * compatibility. It's just making the current lexeme available in
303              * yytext and '\0' terminating the lexeme. The '\0' is removed at 2.
304              * The problem is that you have to read the next lookahead symbol
305              * before you can reduce by the production that had the previous
306              * symbol at its far right. Note that, since Production 0 has the
307              * goal symbol on its left-hand side, a reduce by 0 is an accept
308              * action. Also note that ii_ptext()[ii_plength()] is used at (2)
309              * rather than yytext[yyleng] because the user might have modified
310              * yytext or yyleng in an action.
311             *
312             * Rather than pushing junk as the $$=$1 action on an epsilon
313             * production, the old tos item is duplicated in this situation.
314             */
315
316             act_num = -act_num ;
317             Yy_rhslen = Yy_reduce[ act_num ];
318             Yy_val = Yy_vsp[ Yy_rhslen ? Yy_rhslen-1 : 0 ];    /* $$ = $1 */
319
320             yylineno = ii_lineno();                                /* (1) */
321             yytext = ii_ptext();                                 ;
322             tchar = yytext[ yyleng = ii_plength() ];
323             yytext[yyleng] = '\0' ;
324
325             YYD( yycomment("Reduce by (%d) %s->%s\n", act_num, \
326                                     Yy_slhs[act_num], Yy_srhs[act_num]); )
327
328             if( errcode = yy_act( act_num, Yy_vsp ) )
329                 return errcode;
330

```



Listing 5.14. continued...

```

331         ii_ptext() [ ii_plength() ] = tchar;                      /* (2)      */
332
333     if( act_num == YY_IS_ACCEPT )
334         break;
335     else
336         YY_reduce( act_num, Yy_rhslen );
337     }
338 }
339 YYD( yycomment( "Accept\n" );           )
340 YYD( yy_quit_debug();                  )
341
342 YYACCEPT;
343 }
```

Reduce subroutine,
`yy_reduce()`.

Reductions are done in `yy_reduce()`, on line 128 of Listing 5.14. The subroutine pops as many stack items as are on the right-hand side of the production by which the parser's reducing on lines 134 to 136; the goto state is computed on line 138, and is pushed on line 142 (153 if debugging is enabled). Nothing is popped in the case of an ϵ production because amount is zero in this situation. `Yy_val` (which holds `$$`) is initialized to `$1` before `yy_reduce` is called.⁹ Similarly, the action code (in which `Yy_val` could be modified via `$$`) is executed after the initialization but before the `yy_reduce()` call. The code on lines 144 to 158 of Listing 5.14 is activated in debugging mode. The actions that I just described are performed here too, but they're done in several steps so that breakpoints and screen updates happen in reasonable places.

Stack initialization.

The stacks are all initialized in `yy_init_stack()` on line 165. This initialization must be done at run time because the parser might be called several times within a program (if, for example, it's being used to parse expressions in some large program rather than for a complete compiler). Initially, 0 is pushed onto the state stack, garbage onto the value stack, and the string `"$"` onto the debugging symbol stack.

Error recovery:
`yy_recover()`.

The next issue is error recovery, which is handled by `yy_recover()` on line 182 of Listing 5.15. This routine is passed the token that caused the error. It prints an error message if its second argument is false, and tries to recover from the error by manipulating the stack. The *panic-mode* error-recovery technique described earlier is used:

- (1) Pop items off the stack until you enter a state that has a legitimate action on the current token (for which there is a nonerror transition on the current token). If you find such a state, you have recovered. Return the input token.
- (2) Otherwise, if there are no such states on the stack, advance the input, restore the stack to its original condition (that is, to the condition it was in when the error occurred) and go to (1).
- (3) If you get to the end of input, you can't recover from the error. Return 0 in this case.

The algorithm is represented in pseudocode in Table 5.14.

The parser itself:
`yyparse()`.

The actual parser, `yyparse()`, begins on line 245 of Listing 5.15. It is a straightforward implementation of the simplified, bottom-up parse algorithm in Listing 5.4 on page 373. The remainder of the output file is the third portion of the input file from Listing 5.5

9. **Yacc** hackers: Don't confuse this variable with `yylval`, discussed earlier. **Yacc** uses `yylval` both for `$$` and for the default shift value—**oocs** uses two different variables.

Table 5.14. Bottom-Up, Error-Recovery Algorithm

```

recover( tok )
{
    remember the_real_stack_pointer ;

    do {
        sp = the_real_stack_pointer ;

        while( stack not empty  &&  no_legal_transition( sp, tok ) )
            --sp;

        if( the stack is not empty )
        {
            the_real_stack_pointer = sp;
            return tok;
        }
    } while( (tok = advance( )) ≠ end_of_input );

    return 0;
}

```

on page 383. It is not reproduced here.

5.11 Implementing an LALR(1) Parser Generator— Occs Internals

This section discusses the source code for `occ`, a yacc-like program that translates an augmented, attributed grammar into the C source code for a parser. If you still haven't done it, you must read Appendix E, which is a user's manual for `occ`, before continuing. (I really mean it this time.)

5.11.1 Modifying the Symbol Table for LALR(1) Grammars

Most of the important differences between `LLama` and `occ` are in the `do_file()` subroutine (Listing 4.36, page 327, line 272). The `OX()` macros conditionally compile the code that is unique to `occ`. [`LL()` is used for `LLama`-specific code.] There are other trivial changes scattered throughout the program, but the important ones are in `do_file()`.

The first order of business is modifying the grammar somewhat. `LLama`'s top-down parser can handle actions imbedded in the middle of a production because it just puts them onto the parse stack like any other symbol. The bottom-up parser used by `occ` has to do the actions as part of a reduction, however. Consequently, it has to handle actions differently. The `patch()` subroutine (on line 29 of Listing 5.15) modifies the grammar to be suitable for `occ`, and at the same time removes all actions from the PRODUCTIONs themselves. Actions are output as cases in the `switch` statement in the `yy_act()` subroutine (discussed earlier on page 388). The output `case` statements reference the production number with which the action is associated. The action code itself can safely be discarded once it is output because `occ` no longer needs it.

Modifying the grammar
for `occ`: 'do_patch()'

Listing 5.15. *yypatch.c*—Modify Grammar for Use in Occs

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <malloc.h>
4
5 #include <tools/debug.h>
6 #include <tools/set.h>
7 #include <tools/hash.h>
8 #include <tools/compiler.h>
9 #include <tools/l.h>
10 #include "parser.h"
11
12 /*-----*/
13
14 void patch P(( void )); /* public */
15
16 void dopatch P(( SYMBOL *sym )); /* local */
17 void print_one_case P(( int case_val, unsigned char *action, \
18                         int rhs_size, int lineno, struct _prod_ *prod ));
19
20 /*-----*/
21
22 PRIVATE int Last_real_nonterm ; /* This is the number of the last */
23                                         /* nonterminal to appear in the input */
24                                         /* grammar [as compared to the ones */
25                                         /* that patch() creates]. */
26
27 /*-----*/
28
29 PUBLIC void patch()
30 {
31     /* This subroutine does several things:
32      *
33      *     * It modifies the symbol table as described in the text.
34      *     * It prints the action subroutine and deletes the memory associated
35      *         with the actions.
36      *
37      * This is not a particularly good subroutine from a structured programming
38      * perspective because it does two very different things at the same time.
39      * You save a lot of code by combining operations, however.
40      */
41
42     void dopatch();
43
44     static char *top[] =
45     {
46         "",
47         "yy_act( yypnum, yyvsp )",
48         "int          yypnum;           /* production number */",
49         "YYSTYPE      *yyvsp;           /* value-stack pointer */",
50         "{",
51         "    /* This subroutine holds all the actions in the original input",
52         "     * specification. It normally returns 0, but if any of your",
53         "     * actions return a non-zero number, then the parser halts",
54         "     * immediately, returning that nonzero number to the calling",
55         "     * subroutine.",
56         "    */",
57         "",
58         "    switch( yypnum )",
59         "    {",

```



Listing 5.15. continued...

```

60         NULL
61     };
62
63     static char *bot[] =
64     {
65         "",           /* */
66         "#ifdef YYDEBUG",
67         "        default: yycomment(\"Production %d: no action\\n\", yypnum);",
68         "        break;",
69         "#endif",
70         "    },",
71         "",           /* */
72         "    return 0;",
73         "};",
74     NULL
75     };
76
77     Last_real_nonterm = Cur_nonterm;
78
79     if( Make_actions )
80         printv( Output, top );
81
82     ptab( Symtab, dopatch, NULL, 0 );
83
84     if( Make_actions )
85         printv( Output, bot );
86 }
87
88 /*-----*/
89
90 PRIVATE void dopatch(sym)
91 SYMBOL *sym;
92 {
93     PRODUCTION *prod;      /* Current right-hand side of sym */          */
94     SYMBOL    **p2;        /* General-purpose pointer */                  */
95     SYMBOL    **pp;        /* Pointer to one symbol on rhs */             */
96     SYMBOL    *cur;        /* Current element of right-hand side */       */
97
98     if( !ISNONTERM(sym) || sym->val > Last_real_nonterm )
99     {
100         /* If the current symbol isn't a nonterminal, or if it is a nonterminal
101            * that used to be an action (one that we just transformed), ignore it.
102            */
103         return;
104     }
105
106     for( prod = sym->productions; prod ; prod = prod->next )
107     {
108         if( prod->rhs_len == 0 )
109             continue;
110
111         pp = prod->rhs + (prod->rhs_len - 1);
112         cur = *pp;
113
114         if( ISACT(cur) )                      /* Check rightmost symbol */
115         {
116             print_one_case( prod->num, cur->string, --(prod->rhs_len),
117                               cur->lineno, prod );
118

```



Listing 5.15. continued...

```

119         delsym ( (HASH_TAB*) Symtab, (BUCKET*) cur );
120         free   ( cur->string );
121         freesym ( cur );
122         *pp-- = NULL;
123     }
124     /* cur is no longer valid because of the --pp above. */
125     /* Count the number of nonactions in the right-hand */
126     /* side. */                                                 */
127
128     for( ; pp >= prod->rhs; --pp )
129     {
130         cur = *pp;
131
132         if( !ISACT(cur) )
133             continue;
134
135         if( Cur_nonterm >= MAXNONTERM )
136             error(1,"Too many nonterminals & actions (%d max)\n", MAXTERM);
137         else
138         {
139             /* Transform the action into a nonterminal. */
140
141             Terms[ cur->val = ++Cur_nonterm ] = cur ;
142
143             cur->productions = (PRODUCTION*) malloc( sizeof(PRODUCTION) );
144             if( !cur->productions )
145                 error(1, "INTERNAL [dopatch]: Out of memory\n");
146
147             print_one_case( Num_productions, /* Case value to use.      */
148                            cur->string,        /* Source code.          */
149                            pp - prod->rhs,    /* # symbols to left of act.*/
150                            cur->lineno,        /* Input line # of code. */
151                            prod
152                         );
153
154             /* Once the case is printed, the string argument can be freed.*/
155
156             free( cur->string );
157             cur->string           = NULL;
158             cur->productions->num   = Num_productions++;
159             cur->productions->lhs    = cur ;
160             cur->productions->rhs_len = 0;
161             cur->productions->rhs[0] = NULL;
162             cur->productions->next   = NULL;
163             cur->productions->prec    = 0;
164
165             /* Since the new production goes to epsilon and nothing else,
166              * FIRST(new) == { epsilon }. Don't bother to refigure the
167              * follow sets because they won't be used in the LALR(1) state-
168              * machine routines [If you really want them, call follow()
169              * again.]
170             */
171
172             cur->first = newset();
173             ADD( cur->first, EPSILON );
174         }
175     }
176 }
177 }
178 }
```

Listing 5.15. continued...

```

179 PRIVATE void print_one_case( case_val, action, rhs_size, lineno, prod )
180 int case_val; /* Numeric value attached to case itself.*/
181 unsigned char *action; /* Source Code to execute in case. */
182 int rhs_size; /* Number of symbols on right-hand side. */
183 int lineno; /* input line number (for #lines). */
184 PRODUCTION *prod; /* Pointer to right-hand side. */
185 {
186     /* Print out one action as a case statement. All $-specifiers are mapped
187     * to references to the value stack: $$ becomes Yy_vsp[0], $1 becomes
188     * Yy_vsp[-1], etc. The rhs_size argument is used for this purpose.
189     * [see do_dollar() in yydollar.c for details].
190     */
191
192     int num, i;
193     char *do_dollar(); /* source found in yydollar.c */
194     extern char *production_str(); /* source found in acts.c */
195     char fname[40], *fp; /* place to assemble $<fname>1 */
196
197     if( !Make_actions )
198         return;
199
200     output("\n      case %d: /* %s */\n\n\t", case_val, production_str(prod));
201
202     if( !No_lines )
203         output("#line %d \"%s\"\n\t", lineno, Input_file_name );
204
205     while( *action )
206     {
207         if( *action != '$' )
208             output( "%c", *action++ );
209         else
210         {
211             /* Skip the attribute reference. The if statement handles $$ the
212             * else clause handles the two forms: $N and $-N, where N is a
213             * decimal number. When we hit the do_dollar call (in the output()
214             * call), "num" holds the number associated with N, or DOLLAR_DOLLAR
215             * in the case of $$.
216             */
217
218             if( *++action != '<' )
219                 *fname = '\0';
220             else
221             {
222                 ++action; /* skip the < */
223                 fp = fname;
224
225                 for(i=sizeof(fname); --i>0 && *action && *action != '>'; )
226                     *fp++ = *action++;
227
228                 *fp = '\0';
229                 if( *action == '>' )
230                     ++action;
231             }
232
233             if( *action == '$' )
234             {
235                 num = DOLLAR_DOLLAR;
236                 ++action;
237             }

```



Listing 5.15. continued...

```

238         else
239         {
240             num = atoi( action );
241             if( *action == '-' )
242                 ++action ;
243             while( isdigit(*action) )
244                 ++action ;
245         }
246
247         output( "%s", do_dollar( num, rhs_size, lineno, prod, fnat ) );
248     }
249
250     output( "\n          break;\n" );
251 }
```

Imbedded actions.

Those actions that are already at the far right of a production are just output with the production number as the case value (by the **if** statement in `do_patch()` on lines 114 to 123 of Listing 5.15). The action is effectively removed from the production with the `*pp---=NULL` on line 122. The **for** loop that starts on line 128 looks for imbedded actions, and, if any are found, modifies the grammar accordingly. Productions of the form:

```
s : symbols {action();} symbols ;
```

are translated into:

```
s : symbols 001 symbols
;
001 : { action(); }
```

The new nonterminal (named 001, above) is created on lines 141 to 145 (and initialized on lines 156 to 173) by transforming the existing SYMBOL structure (the one that represents the action) into a nonterminal that has an empty right-hand side. The new production number is assigned on line 141. Note that the production in which the action was imbedded doesn't need to be modified. The SYMBOL pointer that used to reference the action now references the nonterminal, but it's the same pointer in both situations.

Dollar-attribute translation: `do_dollar()`.

%union fields.

The `print_one_case()` routine, which prints the code associated with the action, starts on line 179 of Listing 5.15. Dollar attributes (\$\$, \$1, \$-1, and so forth) are handled on lines 209 to 248. The actual work is done by `do_dollar()`, in Listing 5.16. It is passed the attribute number (the 1 in \$1), the number of symbols on the right-hand side of the current production, the input line number (for error messages) and a pointer to the PRODUCTION. It evaluates to a string holding the code necessary to reference the desired attribute (using the translations discussed earlier on Page 388).

Fields that are attached to a %union are also handled here (on lines 25 to 36 and 57 to 75 of Listing 5.16). The code on lines 25 to 36 handles any fields associated with \$\$, the names of which are found in `*prod->lhs->field`. The code on lines 57 to 75 handles other attributes. Since num is the N in \$N, `(prod->rhs)[num-1]->field` gets the field associated with the specified element of the right-hand side. Note that fields can't be used with negative attribute numbers (because `occ`s has no way of knowing what right-hand side corresponds to that field). A warning is printed in this case.

Listing 5.16. *yydollar.c*—Process \$ Attributes

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/set.h>
4 #include <tools/hash.h>
5 #include "parser.h"
6
7 PUBLIC char *do_dollar( num, rhs_size, lineno, prod, fname )
8 int num;           /* The N in $N, DOLLAR_DOLLAR for $$ (DOLLAR_DOLLAR) */
9          /* is defined in parser.h, discussed in Chapter Four. */
10 int rhs_size;      /* Number of symbols on right-hand side, 0 for tail */
11 int lineno;        /* Input line number for error messages */
12 PRODUCTION *prod; /* Only used if rhs_size is >= 0 */
13 char *fname;       /* name in $<name>N */
14 {
15     static char buf[ 128 ];
16     int i, len ;
17
18     if( num == DOLLAR_DOLLAR )                                /* Do $$ */
19     {
20         strcpy( buf, "Yy_val" );
21
22         if( *fname )                                         /* $<name>N */
23             sprintf( buf+6, ".%s", fname );
24
25     else if( fields_active() )
26     {
27         if( *prod->lhs->field )
28             sprintf( buf+6, ".%s", prod->lhs->field );
29     else
30     {
31         error( WARNING, "Line %d: No <field> assigned to $$, "
32                "using default int field\n", lineno );
33
34         sprintf( buf+6, ".%s", DEF_FIELD );
35     }
36 }
37
38 else
39 {
40     if( num < 0 )
41         ++num;
42
43     if( rhs_size < 0 )                                     /* $N is in tail */
44         sprintf( buf, "Yy_vsp[ Yy_rhslen-%d ]", num );
45     else
46     {
47         if( (i = rhs_size - num) < 0 )
48             error( WARNING, "Line %d: Illegal $%d in production\n",
49                      lineno, num );
50     else
51     {
52         len = sprintf( buf, "yyvsp[%d]", i );
53
54         if( *fname )                                         /* $<name>N */
55             sprintf( buf + len, ".%s", fname );
56
57     else if( fields_active() )
58     {

```



Listing 5.16. continued...

```

59             if( num <= 0 )
60             {
61                 error(NONFATAL, "Can't use %%union field with negative"
62                               " attributes. Use $<field>-N\n" );
63             }
64             else if( * (prod->rhs) [num-1]->field )
65             {
66                 sprintf( buf + len, ".%s", (prod->rhs) [num-1]->field );
67             }
68             else
69             {
70                 error( WARNING, "Line %d: No <field> assigned to $%d, "
71                               "using default int field\n",
72                               lineno, num );
73                 sprintf( buf + len, ".%s", DEF_FIELD );
74             }
75         }
76     }
77 }
78
80     return buf;
81 }
```

5.12 Parser-File Generation

The routines in Listing 5.17 and 5.18 handle the `occ`s parser-file generation. They are much like the `LLama` routines with the same names. There are two possible template files, however. The one that we've been looking at (`occ.par`) is used in most situations. If `-a` is specified on the command line, however, the file `occ.act.par` (in Listing 5.19) is used instead. All that's needed here are external declarations that give us access to variables declared in the associated parser file (generated with a `-p` command-line switch).

Listing 5.17. *yycode.c*—Controller Routine for Table Generation

```

1 void    tables()
2 {
3     make_yy_stok();                                /* in stok.c      */
4     make_token_file();                            /* in stok.c      */
5     make_parse_tables();                         /* in yystate.c */
```

5.13 Generating LALR(1) Parse Tables

The only part of `occ`s we've yet to explore is the table-generation subroutines, all concentrated into a single (somewhat large) module called `yystate.c`. The routines create an LR(1) parse table, but before creating a new LR(1) state, the code looks for an existing state with the same LR(0) kernel items as the new one. If such a state exists, lookahead heads are added to the existing state rather than creating a new one.

Listing 5.18. *yydriver.c*—Routines to Create Ocs Output File

```

1 #include <stdio.h>
2 #include <stdlib.h>
3
4 #include <tools/debug.h>
5 #include <tools/set.h>
6 #include <tools/hash.h>
7 #include <tools/compiler.h>
8 #include <tools/l.h>
9 #include "parser.h"
10
11 /*-----*/
12
13 void file_header      P(( void ));           /* public */
14 void code_header       P(( void ));           /* public */
15 void driver            P(( void ));           /* public */
16
17 /*-----*/
18
19 PRIVATE FILE *Driver_file = stderr ;
20
21 /*-----*/
22 * Routines in this file are occs specific. There's a different version of all
23 * these routines in lldriver.c. They MUST be called in the following order:
24 *          file_header()
25 *          code_header()
26 *          driver()
27 */
28 */
29
30 PUBLIC void file_header()
31 {
32     /* This header is printed at the top of the output file, before the
33      * definitions section is processed. Various #defines that you might want
34      * to modify are put here.
35 */
36
37     output( "#include \"%s\"\n\n", TOKEN_FILE );
38
39     if( Public )
40         output( "#define PRIVATE\n" );
41
42     if( Debug )
43         output( "#define YYDEBUG\n" );
44
45     if( Make_actions )
46         output( "#define YYACTION\n" );
47
48     if( Make_parser )
49         output( "#define YYPARSER\n" );
50
51     if( !( Driver_file = driver_1(Output, !No_lines, Template) ) )
52         error( NONFATAL, "%s not found--output file won't compile\n", Template );
53 }
54
55 /*-----*/
56
57 PUBLIC void code_header()
58 {

```



Listing 5.18. continued...

```

59     /* This stuff is output after the definitions section is processed, but
60      * before any tables or the driver is processed.
61      */
62
63     driver_2( Output, !No_lines );
64 }
65
66 /*-----*/
67
68 PUBLIC void     driver()
69 {
70     /* Print out the actual parser by copying llama.par to the output file.
71     */
72
73     if( Make_parser )
74         driver_2( Output, !No_lines );
75
76     fclose( Driver_file );
77 }
```

Listing 5.19. *occs-act.par*—File Header for -a Output

```

1 #ifdef YYDEBUG
2 # define YYD(x) x
3 #else
4 # define YYD(x) /* empty */
5 #endif
6 ^L /* User-supplied code from definitions section goes here */
7 #ifndef YYSTYPE /* Default value stack type */
8 # define YYSTYPE int
9 #endif
10
11 #undef YYD /* Redefine YYD in case YYDEBUG was defined */
12 #ifdef YYDEBUG /* explicitly in the header rather than with */
13 # define YYD(x) x /* a -D on the occs command line. */
14 # define printf yycode /* Make printf() calls go to output window. */
15 #else
16 # define YYD(x) /* empty */
17 #endif
18
19 extern void yycode    ();
20 extern void yydata    ();
21 extern void yybss    ();
22 extern void yycomment ();
23
24 extern YYSTYPE *Yy_vsp; /* Value-stack pointer */
25 extern YYSTYPE Yy_val; /* Must hold $$ after act is performed */
26 extern int     Yy_rhslen; /* number of symbols on RHS of current production */
```

Representing LR(1)
items:, ITEM.

Though this process is straightforward, a lot of code and several data structures have to work in concert to implement it. These data structures start on line 25 of Listing 5.20 with an ITEM declaration, which represents an LR(1) item. There is a certain amount of redundancy here in order to speed up the table generation process. For example, the right_of_dot field points to the symbol to the right of the dot. (Remember, the PRODUCTION itself is stored as an array of SYMBOL pointers. It holds NULL if the dot is at

the far right of the production.) `Occs` could extrapolate this information from the PRODUCTION and dot position (the offset to which is in `dot_posn`) every time it used the ITEM, but it's best, for speed reasons, to do the extrapolation only once since the dot position does not change within a given ITEM. Similarly, the production number could be derived from the PRODUCTION structure, but I've copied it here to avoid an extra level of indirection every time the production number is used. The `RIGHT_OF_DOT` macro on line 35 of Listing 5.20 gets the numeric, tokenized value of the SYMBOL to the right of the dot. It evaluates to zero (which isn't used for any of the input symbols) if the dot is at the far right of the production.

Listing 5.20. `yystate.c`— Universal Constants and LR(1) Items

```

1 #include <stdio.h>
2 #include <malloc.h>
3
4 #include <tools/debug.h>
5 #include <tools/set.h>
6 #include <tools/hash.h>
7 #include <tools/compiler.h>
8 #include <tools/l.h>
9 #include "parser.h"
10 #include "lntok.h"           /* For _EOI_ definition */
11
12 /*****
13                                     /* For statistics only:          */
14 PRIVATE int      Nitems        = 0;   /* number of LR(1) items       */
15 PRIVATE int      Npairs        = 0;   /* # of pairs in output tables */
16 PRIVATE int      Ntab_entries = 0;   /* number of transitions in tables */
17 PRIVATE int      Shift_reduce = 0;   /* number of shift/reduce conflicts */
18 PRIVATE int      Reduce_reduce = 0;  /* number of reduce/reduce conflicts */
19
20 #define MAXSTATE    512    /* Max # of LALR(1) states.      */
21 #define MAXOBUF     256    /* Buffer size for various output routines */
22
23 /*****
24
25 typedef struct _item_          /* LR(1) item:                  */
26 {
27     int            prod_num;    /* production number             */
28     PRODUCTION    *prod;       /* the production itself         */
29     SYMBOL         *right_of_dot; /* symbol to the right of the dot */
30     unsigned char  dot_posn;   /* offset of dot from start of production */
31     SET            *lookaheads;  /* set of lookahead symbols for this item */
32
33 } ITEM;
34
35 #define RIGHT_OF_DOT(p) ( (p)->right_of_dot ? (p)->right_of_dot->val : 0 )

```

The next data structure of interest represents an LR(1) state. The STATE structure is defined, along with a few related constants, in Listing 5.21. Most of the closure items don't have to be stored in the STATE because they can be derived from the kernel items if necessary. The exceptions are items that contain ϵ productions, because these items cause reductions from the current state, not a transition to a new state. The closure items with ϵ productions are effectively part of the kernel, at least from the perspective of the parse table. The set of ϵ items are in their own array because the table-creation code occasionally compares two STATE structures to see if they're equivalent. This equivalence can be determined by comparing the true kernel items only. There's no

Representing an LR(1)
state: STATE.

point in also comparing the ϵ items too, because the same kernel items in both states generate the same closure items—including the ϵ items—in both states. The `kernel_items[]` array is kept sorted to facilitate STATE comparisons. The sort criteria is pretty much arbitrary, but the sorting serves to get the two sets of kernel items to appear in the same order in both STATES. I'll discuss the comparison routine that's used for this purpose in a moment.

Listing 5.21. `yystate.c`—LR(1) States

```

36 #define MAXKERNEL 32 /* Maximum number of kernel items in a state. */
37 #define MAXCLOSE 128 /* Maximum number of closure items in a state (less
38 /* the epsilon productions). */
39 #define MAXEPSILON 8 /* Maximum number of epsilon productions that can be
40 /* in a closure set for any given state. */
41
42 typedef short STATENUM;
43
44 typedef struct _state_ /* LR(1) state */
45 {
46     ITEM *kernel_items [MAXKERNEL]; /* Set of kernel items. */
47     ITEM *epsilon_items [MAXEPIILON]; /* Set of epsilon items. */
48
49     unsigned nkitems : 7; /* # items in kernel_items[]. */
50     unsigned neitems : 7; /* # items in epsilon_items[]. */
51     unsigned closed : 1; /* State has had closure performed. */
52
53     STATENUM num; /* State number (0 is start state). */
54 }
55 } STATE;

```

Internal representation of
the parse table:
Actions, ACT, Goto.

Memory management:
`new()`.

CHUNK.

The next listing (Listing 5.22) contains definitions used to build the internal representation of the actual parse tables. The mechanism used is essentially the same as that used by the output tables, but the transitions are stored as linked lists of structures rather than input-symbol/next-state pairs. The `Actions[]` array (on line 64) is indexed by current state. Each cell of `Actions[]` points at the head of a linked list of ACT structures, one element of which is an input symbol, and the other element of which is the action to take when that symbol is encountered from the current state. The `Goto[]` table (on line 68) is essentially the same.

The table-generation code begins in Listing 5.23 with several routines to manage the parse table just discussed. The allocation routine, `new()` on line 73 allocates space for an ACT structure. Since the space used by these structures need never be freed, `malloc()`, which is both slow and expensive in terms of system resources, need not be called every time a new structure is required; rather, `new()` allocates an array of structures with a single `malloc()` call on line 84. `CHUNK`, **#defined** on line 71, determines the number of structures in the array. `new()` then returns pointers to structures in this array until all array elements have been allocated. Only then does it get another array of structures from `malloc()`. The remainder of the routines in the listing either put new elements in the table, or return pointers to existing elements.

Listing 5.22. *yystate.c*—Internal Representation of the Parse Table

```

56 typedef struct act_or_goto
57 {
58     int sym;                      /* Given this input symbol,           */
59     int do_this;                 /* do this. >0 == shift, <0 == reduce   */
60     struct act_or_goto *next;    /* Pointer to next ACT in the linked list. */
61
62 } ACT;
63 typedef ACT GOTO;             /* GOTO is an alias for ACT */
64 PRIVATE ACT *Actions[MAXSTATE]; /* Array of pointers to the head of the action
65                                * chains. Indexed by state number.
66                                * I'm counting on initialization to NULL here.
67                                */
68 PRIVATE GOTO *Gotos[MAXSTATE]; /* Array of pointers to the head of the goto
69                                * chains.
70                                */

```

Listing 5.23. *yystate.c*—Parse-Table Memory Management

```

71 #define CHUNK      128          /* New() gets this many structures at once */
72
73 PRIVATE void      *new()
74 {
75     /* Return an area of memory that can be used as either an ACT or GOTO.
76     * These objects cannot be freed.
77     */
78
79     static ACT  *eheap  = (ACT *) 0;
80     static ACT  *heap   = (ACT *) 1;
81
82     if( heap >= eheap )
83     {
84         if( !(heap = (ACT *) malloc( sizeof(ACT) * CHUNK) ))
85             error( FATAL, "No memory for action or goto\n" );
86
87         eheap = heap + CHUNK ;
88     }
89     ++Ntab_entries ;
90     return heap++ ;
91 }
92
93 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - */
94
95 PRIVATE ACT      *p_action( state, input_sym )
96 int state, input_sym;
97 {
98     /* Return a pointer to the existing ACT structure representing the indicated
99     * state and input symbol (or NULL if no such symbol exists).
100    */
101
102     ACT *p;
103
104     for( p = Actions[state]; p ; p = p->next )
105         if( p->sym == input_sym )
106             return p;
107

```



Listing 5.23. continued...

```

108     return NULL;
109 }
110
111 /* ----- */
112
113 PRIVATE void add_action( state, input_sym, do_this )
114 int state, input_sym, do_this;
115 {
116     /* Add an element to the action part of the parse table. The cell is
117      * indexed by the state number and input symbol, and holds do_this.
118      */
119
120     ACT *p;
121
122     if( Verbose > 1 )
123         printf("Adding shift or reduce action from state %d on %s\n",
124                state, do_this, Terms[ input_sym ]->name);
125
126     p = (ACT *) new();
127     p->sym = input_sym ;
128     p->do_this = do_this ;
129     p->next = Actions[state];
130     Actions[state] = p;
131 }
132 /* ----- */
133
134 PRIVATE GOTO *p_goto( state, nonterminal )
135 int state, nonterminal;
136 {
137     /* Return a pointer to the existing GOTO structure representing the
138      * indicated state and nonterminal (or NULL if no such symbol exists). The
139      * value used for the nonterminal is the one in the symbol table; it is
140      * adjusted down (so that the smallest nonterminal has the value 0)
141      * before doing the table look up, however.
142      */
143
144     GOTO *p;
145
146     nonterminal = ADJ_VAL( nonterminal );
147
148     for( p = Gotos[ state ] ; p ; p = p->next )
149         if( p->sym == nonterminal )
150             return p;
151
152     return NULL;
153 }
154
155 /* ----- */
156
157 PRIVATE void add_goto( state, nonterminal, go_here )
158 int state, nonterminal, go_here;
159 {
160     /* Add an element to the goto part of the parse table, the cell is indexed
161      * by current state number and nonterminal value, and holds go_here. Note
162      * that the input nonterminal value is the one that appears in the symbol
163      * table. It is adjusted downwards (so that the smallest nonterminal will
164      * have the value 0) before being inserted into the table, however.
165      */
166

```



Listing 5.23. continued...

```

167      GOTO      *p;
168      int       unadjusted;           /* Original value of nonterminal */
169
170      unadjusted = nonterminal;
171      nonterminal = ADJ_VAL( nonterminal );
172
173      if(Verbose > 1)
174          printf("Adding goto from state %d to %d on %s\n",
175                  state, go_here, Terms[unadjusted]->name);
176      p      = (GOTO *) new();
177      p->sym = nonterminal;
178      p->do_this = go_here;
179      p->next = Gotos[state];
180      Gotos[state] = p;
181  }

```

The next listing (Listing 5.24) contains definitions for the data structures used to manage the LR(1) STATES as the state machine is constructed. Before looking at the actual code, consider the logic used to create the machine, summarized in Table 5.15.

State management.

Table 5.15. An Algorithm for Creating an LALR(1) State Machine

Data structures:	
A state data base.	
A list of unclosed states in the state data base.	
Initially:	
Create the start state and put it into both the state data base and the unclosed-state list.	
Create LALR(1) state machine:	
a. for(each <i>state</i> in the list of unclosed states)	
{	
b. Perform LR(1) closure on <i>state</i> , generating a set of new states and associated lookaheads.	
c. for(each <i>new_state</i> generated by the closure operation)	
{	
d. if(an <i>existing_state</i> with the same kernel items as the <i>new_state</i> already exists)	
{	
Merge appropriate next-state transitions from <i>state</i> with <i>existing_state</i> .	
Add lookaheads in <i>new_state</i> to the lookaheads in <i>existing_state</i> .	
if(the previous step added new symbols to the lookahead set AND	
the existing state isn't already in the list of unclosed states)	
Add the existing state to the list of unclosed states.	
}	
else	
{	
Add appropriate next-state transitions from <i>state</i> to <i>new_state</i> .	
Add the <i>new_state</i> to the state data base.	
Add the <i>new_state</i> to the list of unclosed states.	
}	
}	

Two data structures are used by the algorithm to keep track of the states. The states themselves must be stored in a data base of some sort because every time a new state is created, you have to see if a state with the same set of kernel items already exists. (If it

Finished-states list.

Listing 5.24. *yystate.c*—LR-State Memory Management—Data Structures

```

182 PRIVATE HASH_TAB *States      = NULL;    /* LR(1) states          */
183 PRIVATE int Nstates         = 0;        /* Number of states.     */
184
185 #define MAX_UNFINISHED 128
186
187 typedef struct tnode
188 {
189     STATE *state;
190     struct tnode *left, *right;
191 }
192 } TNODE;
193
194 PRIVATE TNODE Heap[ MAX_UNFINISHED ]; /* Source of all TNODES */
195 PRIVATE TNODE *Next_allocate = Heap; /* Ptr to next node to allocate */
196
197 PRIVATE TNODE *Available = NULL; /* Free list of available nodes */
198                                         /* linked list of TNODES. p->left */
199                                         /* is used as the link. */
200 PRIVATE TNODE *Unfinished = NULL; /* Tree of unfinished states. */
201
202 PRIVATE ITEM **State_items; /* Used to pass info to state_cmp */
203 PRIVATE int State_nitems; /* " */
204 PRIVATE int Sort_by_number = 0; /* " */
205
206 #define NEW 0 /* Possible return values from */
207 #define UNCLOSED 1 /* newstate(). */
208 #define CLOSED 2 /* */

```

does exist, the next-state transitions go to the existing state rather than the newly created one.) An unsorted array is inappropriate for holding the list of existing states because the search time is too long. Fortunately, we already have a quite workable data-base manager in the guise of the hash-table functions that we've been using for symbol-table management, and I use those routines here to keep track of the existing states. The States pointer (declared on line 182 of Listing 5.24) points at a hash table of LR(1) states. The maketab() call that initializes the table is done in one of the higher-level routines discussed later in this section.

Unfinished-states list.

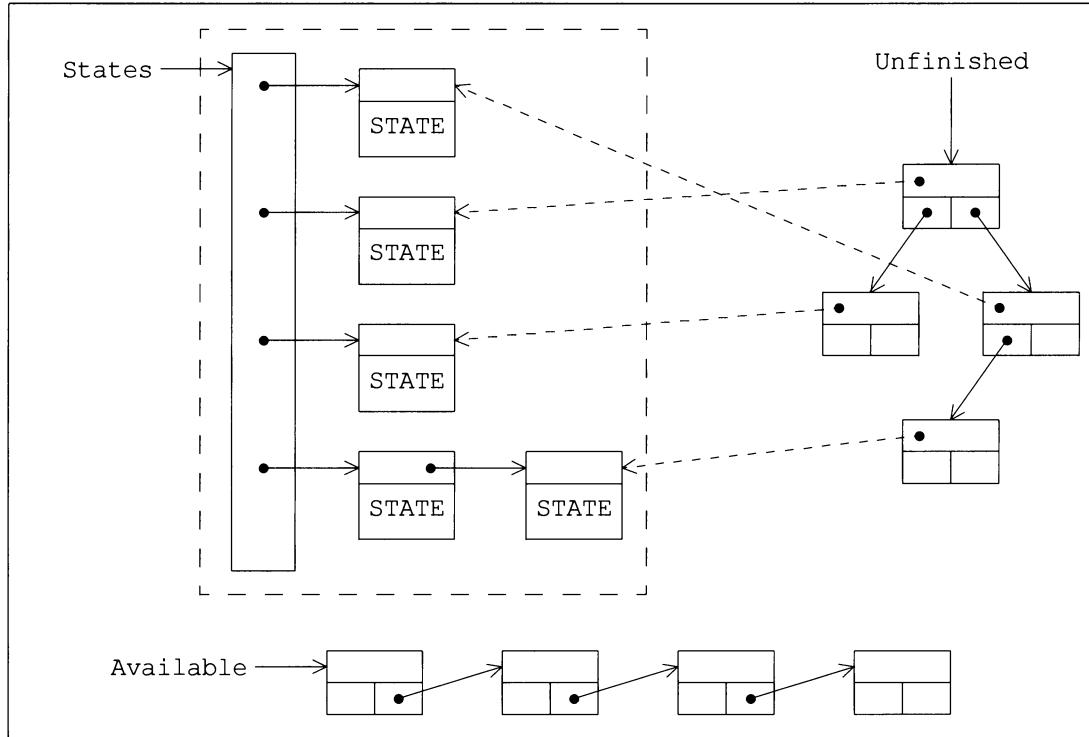
The next data structure keeps track of those states that are created in a closure operation, but have not, themselves, been closed. Since each closure operation creates several new states, a mechanism is needed to keep track of those states on which closure has not yet been performed. Look-up time is an issue here, as well, because step (g) of the algorithm must test to see if the state is already in the list of unclosed states before adding it to the list. Though a hash table could also be used here, I've opted for a binary tree because it's both easy to implement in this application and somewhat more efficient than the general-purpose hash-table functions. It's also easier, in the current application, to remove an arbitrary state from a binary tree than from the hash table—the removed state is always a leaf, here. So, the unclosed-state list is implemented as a binary tree of pointers to STATES in the hash table. The TNODE structure used for the tree nodes is defined on lines 187 to 192 of Listing 5.24. The complete system is pictured in Figure 5.14.

TNODE.

Unfinished-state management: Available.

Next_allocate, Heap.

The Available list at the bottom of Figure 5.14 manages nodes in the tree that have been used at least once. Tree nodes are, initially, managed like the STATES. An array of several of them (called Heap) is defined on line 194 of Listing 5.24. Next_allocate (defined on the next line) is a pointer to the next available node in

Figure 5.14. The State Data Base and Unclosed-State List

Heap[]. Initially, nodes are taken from the Heap[] by incrementing the `Next_allocate` pointer. When a node is removed from the unclosed-state tree, that node is added to the linked list pointed to by `Available`. New nodes are allocated from the linked list if it isn't empty, otherwise a new node is fetched from the Heap[]. This strategy is necessary because `malloc()` and `free()` are too slow.

New states are created and initialized by `newstate`, at the top of Listing 5.25. Note that you don't have to initialize the new node until after you've decided whether or not to put it into the tree. You can get away with this delay because the comparison function used to do the insertion is always called with: `cmp(existing, new)`, where `existing` is the node already in the tree and `new` is the new one. The comparison function doesn't have to actually examine the new node to do the comparison, however. The `ITEM` pointer and count that's passed into `newstate()` can be copied into global variables, and the comparison function can look at those global variables instead of the actual contents of the new node. `State_items` and `State_nitems`, which were declared on lines 202 and 203 of Listing 5.24, are used for this purpose. The delayed initialization gives you a little more speed, because you don't have to do a true initialization unless it's absolutely necessary.

The `add_unfinished()` subroutine on line 267 of Listing 5.25 adds nodes to the unfinished list, but only if the node is not already there. As discussed earlier, a binary tree is used for the list. Note that this is not an ideal data structure because nodes are, at least initially, inserted in ascending order, yielding a worst-case linked list. Things improve once the lookaheads start being added because previously removed states are put back onto the list more or less randomly. This situation could be improved by using a unique random number as an identifier rather than the state number. But, it will most likely take longer to generate this number than it will to chase down the list. Alternately, you could keep the list as an array and do a binary-insertion sort to put in the new

STATE allocation: `newstate()`.

Add to unfinished list:
`add_unfinished()`.

Listing 5.25. *yystate.c*—LR-State Memory Management—Subroutines

```

209 PRIVATE int newstate( items, nitems, statep )
210 ITEM    **items;
211 int     nitems;
212 STATE   **statep;
213 {
214     STATE *state;
215     STATE *existing;
216     int    state_cmp();
217
218     if( nitems > MAXKERNEL )
219         error( FATAL, "Kernel of new state %d too large\n", Nstates );
220
221
222     State_items = items;      /* set up parameters for state_cmp */
223     State_nitems = nitems;    /* and state_hash. */
224
225     if( existing = findsym( States, NULL ) )
226     {
227         /* State exists; by not setting "state" to NULL, we'll recycle */
228         /* the newly allocated state on the next call. */
229
230         *statep = existing;
231         if( Verbose > 1 )
232         {
233             printf("Using existing state (%sclosed): ",
234                               existing->closed ? "" : "un" );
235             pstate_stdout( existing );
236         }
237         return existing->closed ? CLOSED : UNCLOSED ;
238     }
239     else
240     {
241         if( Nstates >= MAXSTATE )
242             error(FATAL, "Too many LALR(1) states\n");
243
244         if( !(state = (STATE *) newsym(sizeof(STATE)) ) )
245             error( FATAL, "Insufficient memory for states\n" );
246
247         memcpy( state->kernel_items, items, nitems * sizeof(ITEM*) );
248         state->nkitems = nitems;
249         state->neitems = 0;
250         state->closed = 0;
251         state->num = Nstates++ ;
252         *statep = state;
253         addsym( States, state );
254
255         if( Verbose > 1 )
256         {
257             printf("Forming new state:");
258             pstate_stdout( state );
259         }
260
261         return NEW;
262     }
263 }
264
265 /*-----*/
266

```



Listing 5.25. continued...

```

267 PRIVATE void add_unfinished( state )
268 STATE *state;
269 {
270     TNODE **parent, *root;
271     int cmp;
272
273     parent = &Unfinished;
274     root = Unfinished;
275     while( root ) /* look for the node in the tree */
276     {
277         if( (cmp = state->num - root->state->num) == 0 )
278             break;
279         else
280         {
281             parent = (cmp < 0) ? &root->left : &root->right ;
282             root = (cmp < 0) ? root->left : root->right ;
283         }
284     }
285
286     if( !root ) /* Node isn't in tree. */
287     {
288         if( Available ) /* Allocate a new node and */
289         {
290             *parent = Available ; /* put it into the tree. */
291             Available = Available->left ; /* Use node from Available */
292         } /* list if possible, otherwise */
293         /* get the node from the Heap. */
294     }
295     else
296     {
297         if( Next_allocate >= &Heap[ MAX_UNFINISHED ] )
298             error(FATAL, "Internal: No memory for unfinished state\n");
299         *parent = Next_allocate++;
300     }
301
302     (*parent)->state = state; /* initialize the node */
303     (*parent)->left = (*parent)->right = NULL;
304 }
305 /*-----*/
306
307 PRIVATE STATE *get_unfinished()
308 {
309     /* Returns a pointer to the next unfinished state and deletes that
310      * state from the unfinished tree. Returns NULL if the tree is empty.
311      */
312
313     TNODE *root;
314     TNODE **parent;
315
316     if( !Unfinished )
317         return NULL;
318
319     parent = &Unfinished; /* find leftmost node */
320     if( root = Unfinished )
321     {
322         while( root->left )
323         {
324             parent = &root->left ;
325             root = root->left ;
326         }
327     }
328 }
```



Listing 5.25. continued...

```

327      }
328
329      *parent    = root->right ;           /* Unlink node from the tree */
330      root->left  = Available;          /* Put it into the free list */
331      Available  = root;
332
333      return root->state ;
334
335

```

elements. Again, the time required to move the tail end of the array is comparable to the time required to chase down the list. Another solution uses a SET of unfinished-state numbers. The problem here is that, towards the end of creating a big table (when things are slower anyway), almost all of the unfinished states will have large state numbers. If the smallest state number is 256 (for example) you'll have to look at 16 empty words in the bit map before finding one with a set bit. Empirical testing showed that using a SET slowed things down by about 5% (which translates to about 10 seconds on a big grammar, with *occs* running on my 8MHz IBM-PC/AT). The *get_unfinished()* subroutine on line 307 of Listing 5.25 gets a node from the unfinished list. It returns 0 when the list is empty.

STATE comparison,
state_cmp().

The next Listing (Listing 5.26) holds a comparison and a hash function that can be used to order the states. The comparison function, *state_cmp()* on line 336 of Listing 5.26, is needed to compare states when searching for a particular state in the array. States are equivalent if they have the same set of kernel items, and equivalence can be determined from the production numbers and dot positions only. The lookaheads don't matter because you're testing for equivalence only to see if the lookaheads for a new LR(1) state can be merged with those of an existing one. Similarly, the actual ordering of states is immaterial as long as the ordering is consistent, so you can do easy things for sort criteria. The rules are:

- (1) The STATE with the most items is larger.
- (2) If the STATES have the same number of items, the two lists of items are compared. Remember, these lists are sorted, so you can do what amounts to a lexicographic comparison, here. The **for** loop on lines 360 to 367 of Listing 5.26 scans the list of items until it finds a mismatched item. If there are no mismatched items, the states are equivalent.
- (3) The items are compared using two criteria: The item with the largest production number is the largest (this test is on line 362). Otherwise, if both items have the same production number, the item with the dot further to the right is the larger (this test is on line 365).

STATE hashing,
state_hash().

ITEM management,
newitem(),
free_item(),
free_recycled_items().
movedot().

ITEM comparison:
item_cmp().

The hash function on line 372 of Listing 5.26 is used to manage the finished-state list. It is much easier to implement than the sort-comparison function. It just sums together the production numbers and dot positions of the items.

The next set of subroutines, in Listing 5.27 takes care of ITEM management. *newitem()* (on line 394) allocates space for (and initializes) a new item, *freeitem()* (on line 421) puts an item on a recycle heap for later allocation by *newitem()*. *free_recycled_items()* (on line 429) cleans up after all processing finishes—it frees memory used for the entire set of items. The *movedot()* function on line 442 moves the dot over one notch and updates associated fields in the ITEM structure.

The *item_cmp()* function on line 365 is used to sort the items within a STATE structure. The primary sort criterion is the numeric value that represents the symbol to

Listing 5.26. *yystate.c*— State Comparison Functions

```

336 PRIVATE int state_cmp( new, tab_node )
337 STATE *new; /* Pointer to new node (ignored) */
338 STATE *tab_node; /* if Sort_by_number is false). */
339 /* Pointer to existing node */
340 {
341     /* Compare two states as described in the text. Return a number representing
342     * the relative weight of the states, or 0 if the states are equivalent.
343     */
344
345     ITEM **tab_item; /* Array of items for existing state */
346     ITEM **item; /* Array of items for new state */
347     int nitem; /* Size of */
348     int cmp;
349
350     if( Sort_by_number )
351         return( new->num - tab_node->num );
352
353     if( cmp = State_nitems - tab_node->nkitems ) /* state with largest */
354         return cmp; /* number of items is */
355     /* larger. */
356     nitem = State_nitems;
357     item = State_items;
358     tab_item = tab_node->kernel_items;
359
360     for( --nitem >= 0 ; ++tab_item, ++item )
361     {
362         if( cmp = (*item)->prod_num - (*tab_item)->prod_num )
363             return cmp;
364
365         if( cmp = (*item)->dot_posn - (*tab_item)->dot_posn )
366             return cmp;
367     }
368
369     return 0; /* States are equivalent */
370 }
371
372 PRIVATE int state_hash( sym )
373 STATE *sym; /* ignored */
374 {
375     /* Hash function for STATES. Sum together production numbers and dot
376     * positions of the kernel items.
377     */
378
379     ITEM **items; /* Array of items for new state */
380     int nitems; /* Size of */
381     int total;
382
383     items = State_items;
384     nitems = State_nitems;
385     total = 0;
386
387     for( --nitems >= 0 ; ++items )
388         total += (*items)->prod_num + (*items)->dot_posn;
389
390     return total;
391 }
```

the right of the dot. This means that ϵ productions float to the top of the list, followed by productions with terminals to the right of the dot, followed by productions with nonterminals to the right of the dot. The ordering is handy when you're partitioning the closure items to create new state kernels. Items with the same symbol to the right of the dot will be adjacent in the sorted list. The second sort criteria (if the same symbol is to the right of the dot in both items) is the production number, and the tertiary criteria is the dot position.

Listing 5.27. *yystate.c*—ITEM Management

```

392 ITEM *Recycled_items = NULL;
393
394 PRIVATE ITEM *newitem( production )
395 PRODUCTION *production;
396 {
397     ITEM *item;
398
399     if( Recycled_items )
400     {
401         item = Recycled_items ;
402         Recycled_items = (ITEM *) Recycled_items->prod;
403         CLEAR( item->lookaheads );
404     }
405     else
406     {
407         if( !(item = (ITEM *) malloc( sizeof(ITEM) ) ) )
408             error( FATAL, "Insufficient memory for all LR(1) items\n" );
409
410         item->lookaheads = newset() ;
411     }
412
413     ++Nitems;
414     item->prod      = production ;
415     item->prod_num   = production->num ;
416     item->dot_posn    = 0;
417     item->right_of_dot = production->rhs[0] ;
418     return item;
419 }
420
421 PRIVATE void freeitem( item )
422 ITEM *item;
423 {
424     --Nitems;
425     item->prod = (PRODUCTION *) Recycled_items ;
426     Recycled_items = item;
427 }
428
429 PRIVATE void free_recycled_items()
430 {
431     /* empty the recycling heap, freeing all memory used by items there */
432
433     ITEM *p;
434
435     while( p = Recycled_items )
436     {
437         Recycled_items = (ITEM *) Recycled_items->prod ;
438         free(p);
439     }
440 }
441

```



Listing 5.27. continued...

```

442 PRIVATE movedot( item )
443 ITEM *item;
444 {
445     /* Moves the dot one position to the right and updates the right_of_dot
446     * symbol.
447     */
448
449     item->right_of_dot = ( item->prod->rhs )[ ++item->dot_posn ] ;
450 }
451
452 PRIVATE int item_cmp( item1p, item2p )
453 ITEM **item1p, **item2p ;
454 {
455     /* Return the relative weight of two items, 0 if they're equivalent. */
456
457     int rval;
458     ITEM *item1 = *item1p;
459     ITEM *item2 = *item2p;
460
461     if( !(rval = RIGHT_OF_DOT(item1) - RIGHT_OF_DOT(item2)) )
462         if( !(rval = item1->prod_num - item2->prod_num ) )
463             return item1->dot_posn - item2->dot_posn ;
464
465     return rval;
466 }
```

The next three listings contain subroutines that make the LR(1) parse table using the logic discussed previously on page 415. The remainder of the chapter is boring stuff that does things like print out the tables. Rather than discuss the routines here, I have commented them copiously so that you don't have to be flipping back and forth between the listing and the text.

Making the parse tables.

Listing 5.28. yystate.c— High-Level, Table-Generation Function

```

467 PUBLIC void make_parse_tables()
468 {
469     /* Prints an LALR(1) transition matrix for the grammar currently
470     * represented in the symbol table.
471     */
472
473     ITEM *item;
474     STATE *state;
475     PRODUCTION *start_prod;
476     void mkstates();
477     int state_cmp();
478     FILE *fp, *old_output;
479
480     /* Make data structures used to produce the table, and create an initial
481     * LR(1) item containing the start production and the end-of-input marker
482     * as a lookahead symbol.
483     */
484
485     States = maketab( 257, state_hash, state_cmp );
486
487     if( !Goal_symbol )
488         error(FATAL, "No goal symbol.\n" );
489 }
```



Listing 5.28. continued...

```

490     start_prod = Goal_symbol->productions ;
491
492     if( start_prod->next )
493         error(FATAL, "Start symbol must have only one right-hand side.\n" );
494
495     item = newitem( start_prod );           /* Make item for start production */
496     ADD( item->lookaheads, _EOI_ );        /* FOLLOW(S) = {} */
497
498     newstate( &item, 1, &state );
499
500     if( lr( state ) )                  /* Add shifts and gotoes to the table */
501     {
502         if( Verbose )
503             printf("Adding reductions:\n");
504
505         reductions();                   /* add the reductions */
506
507         if( Verbose )
508             printf("Creating tables:\n" );
509
510         if( !Make_yyoutab )           /* Tables go in yyout.c */
511         {
512             print_tab( Actions, "Yya", "Yy_action", 1 );
513             print_tab( Gotos , "Yyg", "Yy_goto" , 1 );
514         }
515     else
516     {
517         if( !(fp = fopen(TAB_FILE, "w")) )
518         {
519             error( NONFATAL, "Can't open " TAB_FILE ", ignoring -T\n" );
520
521             print_tab( Actions, "Yya", "Yy_action", 1 );
522             print_tab( Gotos , "Yyg", "Yy_goto" , 1 );
523         }
524     else
525     {
526         output ( "extern YY_TTYPE *Yy_action[]; /* in yyoutab.c */\n" );
527         output ( "extern YY_TTYPE *Yy_goto[]; /* in yyoutab.c */\n" );
528
529         old_output = Output;
530         Output      = fp;
531
532         fprintf(fp, "#include <stdio.h>\n" );
533         fprintf(fp, "typedef short YY_TTYPE;\n" );
534         fprintf(fp, "#define YYPRIVATE %s\n",
535                         Public ? /* empty */ : "static" );
536
537         print_tab( Actions, "Yya", "Yy_action", 0 );
538         print_tab( Gotos , "Yyg", "Yy_goto" , 0 );
539
540         fclose( fp );
541         Output = old_output;
542     }
543 }
544     print_reductions();
545 }
546 }
```

Listing 5.29. *yystate.c*—Table-Generation Functions

```

547 PRIVATE int lr( cur_state )
548 STATE *cur_state;
549 {
550     /* Make LALR(1) state machine. The shifts and gotos are done here, the
551     * reductions are done elsewhere. Return the number of states.
552     */
553
554     ITEM **p;
555     ITEM **first_item;
556     ITEM *closure_items[ MAXCLOSE ];
557     STATE *next;           /* Next state. */ */
558     int isnew;             /* Next state is a new state. */ */
559     int nclose;            /* Number of items in closure_items. */ */
560     int nitems;             /* # items with same symbol to right of dot. */ */
561     int val;               /* Value of symbol to right of dot. */ */
562     SYMBOL *sym;            /* Actual symbol to right of dot. */ */
563     int nlr = 0;            /* Nstates + nlr == number of LR(1) states. */ */
564
565     add_unfinished( cur_state );
566
567     while( cur_state = get_unfinished() )
568     {
569         if( Verbose > 1 )
570             printf( "Next pass...working on state %d\n", cur_state->num );
571
572         /* closure() adds normal closure items to closure_items array.
573          * kclosure() adds to that set all items in the kernel that have
574          * outgoing transitions (ie. whose dots aren't at the far
575          * right).
576          * assort() sorts the closure items by the symbol to the right
577          * of the dot. Epsilon transitions will sort to the head of
578          * the list, followed by transitions on nonterminals,
579          * followed by transitions on terminals.
580          * move_eps() moves the epsilon transitions into the closure kernel set.
581         */
582
583         nclose = closure( cur_state, closure_items, MAXCLOSE );
584         nclose = kclosure( cur_state, closure_items, MAXCLOSE, nclose );
585         if( nclose )
586         {
587             assort( closure_items, nclose, sizeof(ITEM*), item_cmp );
588             nitems = move_eps( cur_state, closure_items, nclose );
589             p = closure_items + nitems;
590             nclose -= nitems ;
591
592             if( Verbose > 1 )
593                 pclosure( cur_state, p, nclose );
594         }
595         /* All of the remaining items have at least one symbol to the    */
596         /* right of the dot. */ */
597
598         while( nclose > 0 )
599         {
600             first_item = p ;
601             sym = (*first_item)->right_of_dot ;
602             val = sym->val ;
603

```



Listing 5.29. continued...

```

604             /* Collect all items with the same symbol to the right of the dot.
605              * On exiting the loop, nitems will hold the number of these items
606              * and p will point at the first nonmatching item. Finally nclose is
607              * decremented by nitems. items = 0 ;
608              */
609
610             nitems = 0 ;
611             do
612             {
613                 movedot( *p++ );
614                 ++nitems;
615
616             } while( --nclose > 0 && RIGHT_OF_DOT(*p) == val );
617
618             /* (1) newstate() gets the next state. It returns NEW if the state
619                * didn't exist previously, CLOSED if LR(0) closure has been
620                * performed on the state, UNCLOSED otherwise.
621                * (2) add a transition from the current state to the next state.
622                * (3) If it's a brand-new state, add it to the unfinished list.
623                * (4) otherwise merge the lookaheads created by the current closure
624                * operation with the ones already in the state.
625                * (5) If the merge operation added lookaheads to the existing set,
626                * add it to the unfinished list.
627             */
628
629             isnew = newstate( first_item, nitems, &next );           /* 1 */
630
631             if( !cur_state->closed )
632             {
633                 if( ISTERM( sym ) )                                /* 2 */
634                     add_action( cur_state->num, val, next->num );
635                 else
636                     add_goto    ( cur_state->num, val, next->num );
637             }
638
639             if( isnew == NEW )
640                 add_unfinished( next );                         /* 3 */
641             else
642             {                                                 /* 4 */
643                 if( merge_lookaheads( next->kernel_items, first_item, nitems))
644                 {
645                     add_unfinished( next );                      /* 5 */
646                     ++nlr;
647                 }
648                 while( --nitems >= 0 )
649                     freeitem( *first_item++ );
650             }
651
652             fprintf( stderr, "\rLR:-3d LALR:-3d", Nstates + nlr, Nstates );
653         }
654         cur_state->closed = 1;
655     }
656
657     free_recycled_items();
658
659     if( Verbose )
660         fprintf(stderr, " states, %d items, %d shift and goto transitions\n",
661                           Nitems, Ntab_entries );
662
663     return Nstates;
}

```



Listing 5.29. continued...

```

664
665 /*-----*/
666
667 PRIVATE int      merge_lookaheads( dst_items, src_items, nitems )
668 ITEM    **src_items;
669 ITEM    **dst_items;
670 int     nitems;
671 {
672     /* This routine is called if newstate has determined that a state having the
673      * specified items already exists. If this is the case, the item list in the
674      * STATE and the current item list will be identical in all respects except
675      * lookaheads. This routine merges the lookaheads of the input items
676      * (src_items) to the items already in the state (dst_items). 0 is returned
677      * if nothing was done (all lookaheads in the new state are already in the
678      * existing state), 1 otherwise. It's an internal error if the items don't
679      * match.
680 */
681
682     int did_something = 0;
683
684     while( --nitems >= 0 )
685     {
686         if(      (*dst_items)->prod      != (*src_items)->prod
687             ||  (*dst_items)->dot_posn != (*src_items)->dot_posn )
688         {
689             error(FATAL, "INTERNAL [merge_lookaheads], item mismatch\n" );
690         }
691
692         if( !subset( (*dst_items)->lookaheads, (*src_items)->lookaheads ) )
693         {
694             ++did_something;
695             UNION( (*dst_items)->lookaheads, (*src_items)->lookaheads );
696         }
697
698         ++dst_items;
699         ++src_items;
700     }
701
702     return did_something;
703 }
704
705 /*-----*/
706
707 PRIVATE int      move_eps( cur_state, closure_items, nclose )
708 STATE   *cur_state;
709 ITEM    **closure_items;
710 {
711     /* Move the epsilon items from the closure_items set to the kernel of the
712      * current state. If epsilon items already exist in the current state,
713      * just merge the lookaheads. Note that, because the closure items were
714      * sorted to partition them, the epsilon productions in the closure_items
715      * set will be in the same order as those already in the kernel. Return
716      * the number of items that were moved.
717 */
718
719     ITEM **eps_items, **p ;
720     int nitems, moved ;
721

```



Listing 5.29. continued...

```

722     eps_items = cur_state->epsilon_items;
723     nitems    = cur_state->neitems ;
724     moved    = 0;
725
726     for( p = closure_items; (*p)->prod->rhs_len == 0 && --nclose >= 0; )
727     {
728         if( ++moved > MAXEPSILON )
729             error(FATAL, "Too many epsilon productions in state %d\n",
730                   cur_state->num );
731         if( nitems )
732             UNION( (*eps_items++)->lookaheads, (*p++)->lookaheads );
733         else
734             *eps_items++ = *p++ ;
735     }
736
737     if( moved )
738         cur_state->neitems = moved ;
739
740     return moved ;
741 }
742
743 /*-----*/
744
745 PRIVATE int kclosure( kernel, closure_items, maxitems, nclose )
746 STATE *kernel;           /* Kernel state to close. */
747 ITEM **closure_items;   /* Array into which closure items are put. */
748 int maxitems;           /* Size of the closure_items[] array. */
749 int nclose;              /* # of items already in set. */
750 {
751     /* Adds to the closure set those items from the kernel that will shift to
752      * new states (ie. the items with dots somewhere other than the far right).
753      */
754
755     int nitems;
756     ITEM *item, **itemp, *citem ;
757
758     closure_items += nclose;           /* Correct for existing items */
759     maxitems -= nclose;
760
761     itemp = kernel->kernel_items ;
762     nitems = kernel->nkitems ;
763
764     while( --nitems >= 0 )
765     {
766         item = *itemp++;
767
768         if( item->right_of_dot )
769         {
770             citem          = newitem( item->prod );
771             citem->prod    = item->prod ;
772             citem->dot_posn = item->dot_posn ;
773             citem->right_of_dot = item->right_of_dot ;
774             citem->lookaheads = dupset( item->lookaheads );
775
776             if( --maxitems < 0 )
777                 error( FATAL, "Too many closure items in state %d\n",
778                         kernel->num );
779             *closure_items++ = citem;
780             ++nclose;
781         }

```



Listing 5.29. continued...

```

782     }
783     return nclose;
784 }
785
786 /*-----*/
787
788 PRIVATE int closure( kernel, closure_items, maxitems )
789 STATE *kernel;           /* Kernel state to close. */
790 ITEM  *closure_items[];   /* Array into which closure items are put */
791 int  maxitems;          /* Size of the closure_items[] array. */
792 {
793     /* Do LR(1) closure on the kernel items array in the input STATE. When
794      * finished, closure_items[] will hold the new items. The logic is:
795      *
796      * (1) for( each kernel item )
797          *      do LR(1) closure on that item.
798      * (2) while( items were added in the previous step or are added below )
799          *      do LR(1) closure on the items that were added.
800
801 */
802
803     int i ;
804     int nclose      = 0 ;           /* Number of closure items */
805     int did_something = 0 ;
806     ITEM **p          = kernel->kernel_items ;
807
808     for( i = kernel->nkitems; --i >= 0 ; )                  /* (1) */
809     {
810         did_something |= do_close( *p++, closure_items, &nclose, &maxitems );
811     }
812
813     while( did_something )                                /* (2) */
814     {
815         did_something = 0;
816         p = closure_items;
817         for( i = nclose ; --i >= 0 ; )
818             did_something |= do_close( *p++, closure_items, &nclose, &maxitems );
819     }
820
821     return nclose;
822 }
823 /*-----*/
824
825 PRIVATE int do_close( item, closure_items, nitems, maxitems )
826 ITEM  *item;
827 ITEM  *closure_items[]; /* (output) Array of items added by closure process */
828 int   *nitems;        /* (input) # of items currently in closure_items[] */
829           /* (output) # of items in closure_items after */
830           /* processing */
831 int   *maxitems;      /* (input) max # of items that can be added */
832           /* (output) input adjusted for newly added items */
833 {
834     /* Workhorse function used by closure(). Performs LR(1) closure on the
835      * input item ([A->b.Cd, e] add [C->x, FIRST(de)]). The new items are added
836      * to the closure_items[] array and *nitems and *maxitems are modified to
837      * reflect the number of items in the closure set. Return 1 if do_close()
838      * did anything, 0 if no items were added (as will be the case if the dot
839      * is at the far right of the production or the symbol to the right of the
840      * dot is a terminal).
841
842 */

```



Listing 5.29. continued...

```

842
843     int      did_something = 0;
844     int      rhs_is_nullable;
845     PRODUCTION *prod;
846     ITEM      *close_item;
847     SET       *closure_set;
848     SYMBOL    **symp;
849
850     if( !item->right_of_dot )
851         return 0;
852
853     if( !ISNONTERM( item->right_of_dot ) )
854         return 0;
855
856     closure_set = newset();
857
858     /* The symbol to the right of the dot is a nonterminal. Do the following:
859      *
860      *(1) for( every production attached to that nonterminal )
861      *      if( the current production is not already in the set of
862      *                      closure items)
863      *      add it;
864      *      if( the d in [A->b.Cd, e] doesn't exist )
865      *          add e to the lookaheads in the closure production.
866      *      else
867      *          The d in [A->b.Cd, e] does exist, compute FIRST(de) and add
868      *          it to the lookaheads for the current item if necessary.
869      */
870
871
872     for( prod = item->right_of_dot->productions; prod ; prod = prod->next )
873     {
874
875         if( !(close_item = in_closure_items(prod, closure_items, *nitems)))
876         {
877             if( --(*maxitems) <= 0 )
878                 error(FATAL, "LR(1) Closure set too large\n");
879
880             closure_items[ (*nitems)++ ] = close_item = newitem( prod );
881             ++did_something;
882         }
883
884         if( !*(symp = & ( item->prod->rhs [ item->dot_posn + 1 ])) ) /* (4) */
885         {
886             did_something |= add_lookahead( close_item->lookaheads, /* (5) */
887                                           item->lookaheads );
888         }
889         else
890         {
891             truncate( closure_set ); /* (6) */
892
893             rhs_is_nullable = first_rhs( closure_set, symp,
894                                         item->prod->rhs_len - item->dot_posn - 1 );
895
896             REMOVE( closure_set, EPSILON );
897
898             if( rhs_is_nullable )
899                 UNION( closure_set, item->lookaheads );
900

```



Listing 5.29. continued...

```

901             did_something |= add_lookahead(close_item->lookaheads, closure_set);
902         }
903     }
904
905     delset( closure_set );
906     return did_something;
907 }
908
909 /*-----*/
910
911 PRIVATE ITEM      *in_closure_items(production, closure_item, nitems)
912 ITEM           **closure_item;
913 PRODUCTION *production;
914 {
915     /* If the indicated production is in the closure_items already, return a
916      * pointer to the existing item, otherwise return NULL.
917      */
918
919     for(; --nitems >= 0 ; ++closure_item )
920         if( (*closure_item)->prod == production )
921             return *closure_item;
922
923     return NULL;
924 }
925
926 /*-----*/
927
928 PRIVATE int      add_lookahead( dst, src )
929 SET      *dst, *src;
930 {
931     /* Merge the lookaheads in the src and dst sets. If the original src
932      * set was empty, or if it was already a subset of the destination set,
933      * return 0, otherwise return 1.
934      */
935
936     if( !IS_EMPTY( src ) && !subset( dst, src ) )
937     {
938         UNION( dst, src );
939         return 1;
940     }
941
942     return 0;
943 }
```

Listing 5.30. yystate.c— Adding Reductions to Tables

```

944 PRIVATE void      reductions()
945 {
946     /* Do the reductions. If there's memory, sort the table by state number */
947     /* first so that yyout.doc will look nice. */
948
949     void addreductions();                                /* below */
950
951     Sort_by_number = 1;
952     if( !ptab( States, addreductions, NULL, 1 ) )
953         ptab( States, addreductions, NULL, 0 );
954 }
```



Listing 5.30. continued...

```

955
956 /*-----*/
957
958 PRIVATE void addreductions( state, junk )
959 STATE *state;
960 void *junk;
961 {
962     /* This routine is called for each state. It adds the reductions using the
963      * disambiguating rules described in the text, and then prints the state to
964      * yyout.doc if Verbose is true. I don't like the idea of doing two things
965      * at once, but it makes for nicer output because the error messages will
966      * be next to the state that caused the error.
967 */
968
969     int i;
970     ITEM **item_p;
971
972     for( i = state->nkitems, item_p = state->kernel_items; --i>=0; ++item_p )
973         reduce_one_item( state, *item_p );
974
975     for( i = state->neitems, item_p = state->epsilon_items; --i>=0; ++item_p )
976         reduce_one_item( state, *item_p );
977
978     if( Verbose )
979     {
980         pstate( state );
981
982         if( state->num % 10 == 0 )
983             fprintf( stderr, "%d\r", state->num );
984     }
985 }
986
987 /*-----*/
988
989 PRIVATE void reduce_one_item( state, item )
990 ITEM *item; /* Reduce on this item */
991 STATE *state; /* from this state */
992 {
993     int token; /* Current lookahead */
994     int pprec; /* Precedence of production */
995     int tprec; /* Precedence of token */
996     int assoc; /* Associativity of token */
997     int reduce_by; /* True if conflict can be resolved */
998     int resolved; /* */
999     ACT *ap;
1000
1001     if( item->right_of_dot ) /* No reduction required */
1002         return;
1003
1004     pprec = item->prod->prec; /* precedence of entire production */
1005
1006     for( next_member(NULL); (token = next_member(item->lookaheads)) >= 0; )
1007     {
1008         tprec = Precedence[token].level; /* precedence of lookahead */
1009         assoc = Precedence[token].assoc; /* symbol */
1010
1011         if( !(ap = p_action( state->num, token )) ) /* No conflicts */
1012         {
1013             add_action( state->num, token, -(item->prod_num) );
1014         }

```



Listing 5.30. continued...

```

1015     else if( ap->do_this <= 0 )
1016     {
1017         /* Resolve a reduce/reduce conflict in favor of the production */
1018         /* with the smaller number. Print a warning. */ */
1019
1020         ++Reduce_reduce;
1021
1022         reduce_by = min( -(ap->do_this), item->prod_num );
1023         ap->do_this = -reduce_by ;
1024
1025         error( WARNING, "State %2d: reduce/reduce conflict "
1026                 "%d/%d on %s (choose %d).\n",
1027                 state->num,
1028                 -(ap->do_this), item->prod_num ,
1029                 token ? Terms[token]->name: "<_EOI_>",
1030                 reduce_by
1031             );
1032     else
1033     {
1034         if( resolved = (pprec && tprec) )
1035             if( tprec < pprec || (pprec == tprec && assoc != 'r') )
1036                 ap->do_this = -( item->prod_num );
1037
1038         if(Verbose > 1 || !resolved )
1039         {
1040             ++Shift_reduce;
1041             error( WARNING, "State %2d: shift/reduce conflict %s/%d"
1042                     " (choose %s) %s\n",
1043                     state->num,
1044                     Terms[token]->name,
1045                     item->prod_num,
1046                     ap->do_this < 0 ? "reduce" : "shift",
1047                     resolved ? "(resolved)" : ""
1048                 );
1049         }
1050     }
1051 }
1052 }
```

Listing 5.31. yystate.c— Statistics Functions

```

1053 PUBLIC void lr_stats( fp )
1054 FILE    *fp;
1055 {
1056     /* Print out various statistics about the table-making process */
1057
1058     fprintf(fp, "%4d      LALR(1) states\n",           Nstates );
1059     fprintf(fp, "%4d      items\n",                  Nitems );
1060     fprintf(fp, "%4d      nonerror transitions in tables\n", Ntab_entries );
1061     fprintf(fp, "%4d/%-4d unfinished items\n", Next_allocate - Heap,
1062                         MAX_UNFINISHED);
1063     fprintf(fp, "%4d      bytes required for LALR(1) transition matrix\n",
1064             (2 * sizeof(char*) * Nstates) /* index arrays */
1065             + Nstates /* count fields */
1066             + (Npairs * sizeof(short)) /* pairs */
1067             );
1068     fprintf(fp, "\n");
```



Listing 5.31. continued...

```

1069 }
1070
1071 /*-----*/
1072
1073 PUBLIC int lr_conflicts( fp )
1074 FILE    *fp;
1075 {
1076     /* Print out statistics for the inadequate states and return the number of
1077     * conflicts.
1078     */
1079
1080     fprintf(fp, "%4d      shift/reduce conflicts\n", Shift_reduce );
1081     fprintf(fp, "%4d      reduce/reduce conflicts\n", Reduce_reduce );
1082     return Shift_reduce + Reduce_reduce ;
1083 }
```

Listing 5.32. yystate.c—Print Functions

```

1084 #define MAX_TOK_PER_LINE 10
1085 PRIVATE int Tokens_printed;          /* Controls number of lookaheads printed */
1086                                     /* on a single line of yyout.doc. */
1087 PRIVATE void sprint_tok( bp, format, arg )
1088     char    **bp;
1089     char    *format;                  /* not used here, but supplied by pset() */
1090     int     arg;
1091 {
1092     /* Print one nonterminal symbol to a buffer maintained by the
1093     * calling routine and update the calling routine's pointer.
1094     */
1095
1096     if      ( arg == -1      ) *bp += sprintf( *bp, "null " );
1097     else if( arg == -2      ) *bp += sprintf( *bp, "empty " );
1098     else if( arg == _EOI_ ) *bp += sprintf( *bp, "$ " );
1099     else if( arg == EPSILON) *bp += sprintf( *bp, "" );
1100     else                *bp += sprintf( *bp, "%s ", Terms[arg]->name );
1101
1102     if( ++Tokens_printed >= MAX_TOK_PER_LINE )
1103     {
1104         *bp += sprintf(*bp, "\n\t\t");
1105         Tokens_printed = 0;
1106     }
1107 }
1108
1109 PRIVATE char  *stritem( item, lookaheads )
1110 ITEM    *item;
1111 int     lookaheads;
1112 {
1113     /* Return a pointer to a string that holds a representation of an item. The
1114     * lookaheads are printed too if "lookaheads" is true or Verbose is > 1
1115     * (-V was specified on the command line).
1116     */
1117
1118     static char buf[ MAXOBUF * 2 ];
1119     char    *bp;
1120     int     i;
```

Listing 5.32. continued...

```

1122     bp = buf;
1123     bp += sprintf( bp, "%s->", item->prod->lhs->name );
1124
1125     if( item->prod->rhs_len <= 0 )
1126         bp += sprintf( bp, "<epsilon>. " );
1127     else
1128     {
1129         for( i = 0; i < item->prod->rhs_len ; ++i )
1130         {
1131             if( i == item->dot_posn )
1132                 *bp++ = '.';
1133
1134             bp += sprintf(bp, " %s", item->prod->rhs[i]->name );
1135         }
1136
1137         if( i == item->dot_posn )
1138             *bp++ = '.';
1139     }
1140
1141     if( lookaheads || Verbose >1 )
1142     {
1143         bp += sprintf( bp, " (production %d, precedence %d)\n\t\t[",
1144                               item->prod->num, item->prod->prec );
1145         Tokens_printed = 0;
1146         pset( item->lookaheads, sprint_tok, &bp );
1147         *bp++ = ']';
1148     }
1149
1150     if( bp - buf >= (MAXOBUF * 2) )
1151         error(FATAL, "Internal [stritem], buffer overflow\n" );
1152
1153     *bp = '\0';
1154     return buf;
1155 }
1156
1157 /*-----*/
1158
1159 PRIVATE void pstate( state )
1160 STATE *state;
1161 {
1162     /* Print one row of the parse table in human-readable form yyout.doc
1163      * (stderr if -V is specified).
1164      */
1165
1166     int i;
1167     ITEM **item;
1168     ACT *p;
1169
1170     document( "State %d:\n", state->num );
1171
1172     /* ----- */
1173     /* Print the kernel and epsilon items for the current state. */
1174
1175     for( i=state->nkitems, item=state->kernel_items ; --i >= 0 ; ++item )
1176         document("    %s\n", stritem(*item, (*item)->right_of_dot==0));
1177
1178     for( i=state->neitems, item=state->epsilon_items ; --i >= 0 ; ++item )
1179         document( "    %s\n", stritem(*item, 1) );
1180

```



Listing 5.32. continued...

```

1181     document( "\n" );
1182
1183     /* - - - - - - - - - - - - - - - - - - - - - - - - - - - */
1184     /* Print out the next-state transitions, first the actions,   */
1185     /* then the gotos.                                         */
1186
1187     for( i = 0; i < MINTERM + USED_TERMS; ++i )
1188     {
1189         if( p = p_action( state->num, i ) )
1190         {
1191             if( p->do_this == 0 )
1192             {
1193                 if( p->sym == _EOI_ )
1194                     document( "Accept on end of input\n" );
1195                 else
1196                     error( FATAL, "INTERNAL: state %d, Illegal accept",
1197                                         state->num );
1198             }
1199             else if( p->do_this < 0 )
1200                 document( "Reduce by %d on %s\n", -(p->do_this),
1201                                         Terms[p->sym]->name );
1202             else
1203                 document( "Shift to %d on %s\n", p->do_this,
1204                                         Terms[p->sym]->name );
1205         }
1206     }
1207
1208     for( i = MINNONTERM; i < MINNONTERM + USED_NONTERMS; i++ )
1209     {
1210         if( p = p_goto(state->num, i) )
1211             document( "Goto %d on %s\n", p->do_this,
1212                                         Terms[i]->name );
1213     }
1214
1215 PRIVATE void      pstate_stdout( state )
1216 STATE      *state;
1217 {
1218     document_to( stdout );
1219     pstate( state );
1220     document_to( NULL );
1221 }
1222
1223 /*-----*/
1224
1225 PRIVATE pclosure( kernel, closure_items, nitems )
1226 STATE      *kernel;
1227 ITEM      **closure_items;
1228 {
1229     printf( "\n%d items in Closure of ", nitems );
1230     pstate_stdout( kernel );
1231
1232     if( nitems > 0 )
1233     {
1234         printf( "      -----closure items:----\n" );
1235         while( --nitems >= 0 )
1236             printf( "      %s\n", stritem( *closure_items++, 0 ) );
1237     }
1238 }
```

Listing 5.33. *yystate.c*—Routines That Create Auxiliary Tables

```

1239 PRIVATE void      make_yy_lhs( prodtab )
1240 PRODUCTION        **prodtab;
1241 {
1242     static char *text[] =
1243     {
1244         "The Yy_lhs array is used for reductions. It is indexed by production",
1245         "number and holds the associated left-hand side adjusted so that the",
1246         "number can be used as an index into Yy_goto.",,
1247         NULL
1248     };
1249 PRODUCTION *prod;
1250 int          i;
1251
1252 comment ( Output, text );
1253 output ( "YYPRIVATE int Yy_lhs[%d] =\n{\n", Num_productions );
1254
1255 for( i = 0; i < Num_productions; ++i )
1256 {
1257     prod = *prodtab++;
1258     output("\t/* %3d */\t%d", prod->num, ADJ_VAL( prod->lhs->val ) );
1259
1260     if( i != Num_productions-1 )
1261         output(",");
1262
1263     if( i % 3 == 2 || i == Num_productions-1 )           /* use three columns */
1264         output( "\n" );
1265 }
1266 output("};\n");
1267 }
1268
1269 /*-----*/
1270
1271 PRIVATE void      make_yy_reduce( prodtab )
1272 PRODUCTION        **prodtab;
1273 {
1274     static char *text[] =
1275     {
1276         "The Yy_reduce[] array is indexed by production number and holds",
1277         "the number of symbols on the right-hand side of the production",
1278         NULL
1279     };
1280 PRODUCTION *prod;
1281 int          i;
1282
1283 comment ( Output, text );
1284 output ( "YYPRIVATE int Yy_reduce[%d] =\n{\n", Num_productions );
1285
1286 for( i = 0; i < Num_productions; ++i )
1287 {
1288     prod = *prodtab++;
1289     output( "\t/* %3d */\t%d", prod->num, prod->rhs_len );
1290
1291     if( i != Num_productions-1 )
1292         output(",");
1293
1294     if( i % 3 == 2 || i == Num_productions-1 )           /* use three columns */
1295         output( "\n" );
1296 }
1297 output("};\n");

```



Listing 5.33. continued...

```

1298 }
1299
1300 /*-----*/
1301
1302 PRIVATE make_yy_slhs( prodtab )
1303 PRODUCTION    **prodtab;
1304 {
1305     static char *text[] =
1306     {
1307         "Yy_slhs[] is a debugging version of Yy_lhs[]. It is indexed by",
1308         "production number and evaluates to a string representing the",
1309         "left-hand side of the production.",
1310         NULL
1311     };
1312
1313     PRODUCTION *prod;
1314     int          i;
1315
1316     comment ( Output, text );
1317     output   ( "YYPRIVATE char *Yy_slhs[%d] =\n{\n", Num_productions );
1318
1319     for( i = Num_productions; --i >= 0 ; )
1320     {
1321         prod = *prodtab++;
1322         output("\t/* %3d */\t\"%s\"", prod->num, prod->lhs->name );
1323         output( i != 0 ? ",\n" : "\n" );
1324     }
1325     output("};\n");
1326 }
1327
1328 PRIVATE make_yy_srhs( prodtab )
1329 PRODUCTION    **prodtab;
1330 {
1331     static char *text[] =
1332     {
1333         "Yy_srhs[] is also used for debugging. It is indexed by production",
1334         "number and evaluates to a string representing the right-hand side of",
1335         "the production.",
1336         NULL
1337     };
1338
1339     PRODUCTION *prod;
1340     int          i, j;
1341
1342     comment ( Output, text );
1343     output   ( "YYPRIVATE char *Yy_srhs[%d] =\n{\n", Num_productions );
1344
1345     for( i = Num_productions; --i >= 0 ; )
1346     {
1347         prod = *prodtab++;
1348         output("\t/* %3d */\t\"", prod->num );
1349
1350         for( j = 0; j < prod->rhs_len ; ++j )
1351         {
1352             output( "%s", prod->rhs[j]->name );
1353             if( j != prod->rhs_len - 1 )
1354                 outc( ' ' );
1355         }
1356     }

```



Listing 5.33. continued...

```

1357         output( i != 0 ? "\",\\n" : "\""\\n" );
1358     }
1359     output("};\\n");
1360 }
1361
1362 /*-----
1363 * The following routines generate compressed parse tables. There's currently
1364 * no way to do uncompressed tables. The default transition is the error
1365 * transition.
1366 */
1367
1368 PRIVATE void    print_reductions()
1369 {
1370     /* Output the various tables needed to do reductions */
1371
1372     PRODUCTION **prodtab;
1373
1374     if(! (prodtab= (PRODUCTION**) malloc(sizeof(PRODUCTION*) * Num_productions)))
1375         error(FATAL,"Not enough memory to output LALR(1) reduction tables\\n");
1376     else
1377         ptab( Symtab, mkprod, prodtab, 0 );
1378
1379     make_yy_lhs    ( prodtab );
1380     make_yy_reduce ( prodtab );
1381
1382     output("#ifdef YYDEBUG\\n");
1383
1384     make_yy_slhs   ( prodtab );
1385     make_yy_srhs   ( prodtab );
1386
1387     output("#endif\\n");
1388     free( prodtab );
1389 }
1390
1391 /*-----*/
1392
1393 PRIVATE void    mkprod( sym, prodtab )
1394 SYMBOL          *sym;
1395 PRODUCTION      **prodtab;
1396 {
1397     PRODUCTION *p;
1398
1399     if( ISNONTERM(sym) )
1400         for( p = sym->productions ; p ; p = p->next )
1401             prodtab[ p->num ] = p ;
1402 }
1403
1404 /*-----*/
1405
1406 PRIVATE void    print_tab( table, row_name, col_name, make_private )
1407 ACT        **table;
1408 char       *row_name;      /* Name to use for the row arrays           */
1409 char       *col_name;      /* Name to use for the row-pointers array   */
1410 int        make_private;  /* Make index table private (rows always private) */
1411 {
1412     /* Output the action .or goto table. */
1413
1414     int        i, j;
1415     ACT        *ele, **elep; /* table element and pointer to same      */
1416     ACT        *e, **p;          */

```

Listing 5.33. continued...

```

1417     int          count;           /* # of transitions from this state, always >0 */
1418     int          column;
1419     SET          *redundant = newset();           /* Mark redundant rows */
1420
1421     static char *act_text[] =
1422     {
1423         "The Yy_action table is action part of the LALR(1) transition",
1424         "matrix. It's compressed and can be accessed using the yy_next()", 
1425         "subroutine, declared below.",
1426         "",
1427         "           Yya000[]={ 3,   5,3   ,   2,2   ,   1,1   };",
1428         "   state number---+ |   | |",
1429         "   number of pairs in list--+ | |",
1430         "   input symbol (terminal)-----+ |",
1431         "   action-----+",
1432         "",
1433         "   action = yy_next( Yy_action, cur_state, lookahead_symbol );",
1434         "",
1435         "   action < 0   -- Reduce by production n,  n == -action.",
1436         "   action == 0   -- Accept. (ie. Reduce by production 0.)",
1437         "   action > 0   -- Shift to state n,  n == action.",
1438         "   action == YYF -- error.",
1439         NULL
1440     };
1441     static char *goto_text[] =
1442     {
1443         "The Yy_goto table is goto part of the LALR(1) transition matrix",
1444         "It's compressed and can be accessed using the yy_next() subroutine,",
1445         "declared below."
1446         "",
1447         " nonterminal = Yy_lhs[ production number by which we just reduced ]",
1448         "",
1449         "           Yyg000[]={ 3,   5,3   ,   2,2   ,   1,1   };",
1450         "   uncovered state--+ |   | |",
1451         "   number of pairs in list--+ | |",
1452         "   nonterminal-----+ |",
1453         "   goto this state-----+",
1454         ""
1455         "   goto_state = yy_next( Yy_goto, cur_state, nonterminal );",
1456         NULL
1457     };
1458
1459     comment( Output, table == Actions ? act_text : goto_text );
1460
1461     /* -----
1462      * Modify the matrix so that, if a duplicate rows exists, only one
1463      * copy of it is kept around. The extra rows are marked as such by setting
1464      * a bit in the "redundant" set. (The memory used for the chains is just
1465      * discarded.) The redundant table element is made to point at the row
1466      * that it duplicates.
1467      */
1468
1469     for( elep = table, i = 0; i < Nstates ; ++elep, ++i )
1470     {
1471         if( MEMBER( redundant,i ) )
1472             continue;
1473
1474         for( p=elep+1, j=i ; ++j < Nstates ; ++p )
1475         {

```

Listing 5.33. continued...

```

1476     if( MEMBER( redundant, j ) )
1477         continue;
1478
1479     ele = *elep;           /* pointer to template chain      */
1480     e   = *p;             /* chain to compare against template */
1481
1482     if( !e || !ele )     /* either or both strings have no elements */
1483         continue;
1484
1485     for( ; ele && e ; ele=ele->next, e=e->next )
1486         if( (ele->do_this != e->do_this) || (ele->sym != e->sym) )
1487             break;
1488
1489     if( !e && !ele )
1490     {
1491         /* Then the chains are the same. Mark the chain being compared
1492          * as redundant, and modify table[j] to hold a pointer to the
1493          * template pointer.
1494         */
1495
1496         ADD( redundant, j );
1497         table[j] = (ACT *) elep;
1498     }
1499 }
1500
1501
1502 /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - -
1503 * Output the row arrays
1504 */
1505
1506 for( elep = table, i = 0 ; i < Nstates ; ++elep, ++i )
1507 {
1508     if( !*elep || MEMBER(redundant, i) )
1509         continue;
1510
1511         /* Count the number of transitions from this state */
1512         count = 0;
1513         for( ele = *elep ; ele` ; ele = ele->next )
1514             ++count;
1515
1516         output("YYPRIVATE YY_TTYPE %s%03d[]=%2d,",row_name, elep-table, count);
1517         /* } */
1518
1519         column = 0;
1520         for( ele = *elep ; ele ; ele = ele->next )
1521         {
1522             ++Npairs;
1523             output( "%3d,%-4d", ele->sym, ele->do_this );
1524
1525             if( ++column != count )
1526                 outc( ',' );
1527
1528             if( column % 5 == 0 )
1529                 output("\n\t\t\t");
1530
1531             output( "};\n" );
1532
1533 }

```

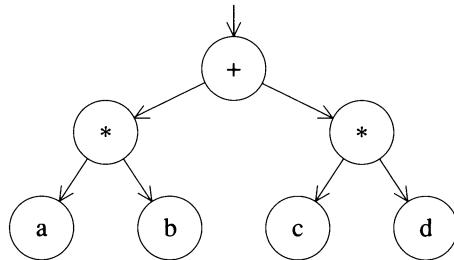
Listing 5.33. continued...

```

1534     /* - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - -
1535     * Output the index array
1536     */
1537
1538     if( make_private )
1539         output( "\nYYPRIVATE YY_TTYPE *%s[%d] =\n", col_name, Nstates );
1540     else
1541         output( "\nYY_TTYPE *%s[%d] =\n", col_name, Nstates );
1542
1543     output( "{\n/* 0 */ " );                                         /* } */
1544
1545     for( elep = table, i = 0 ; i < Nstates ; ++i, ++elep )
1546    {
1547         if( MEMBER(redundant, i) )
1548             output( "%s%03d", row_name, (ACT **)(*elep) - table );
1549         else
1550             output( *elep ? "%s%03d" : " NULL" , row_name, i );
1551
1552         if( i != Nstates-1 )
1553             output( ", " );
1554
1555         if( i==0 || (i % 8)==0 )
1556             output("\n/* %3d */ ", i+1 );
1557     }                                                               /* { */
1558
1559     delset( redundant );                                         /* Mark redundant rows */
1560     output( "\n");\n");
1561 }
```

5.14 Exercises

- 5.1. (a) Write an algorithm that creates a physical syntax tree from a postfix representation of an arithmetic expression. For example, the input ab^*cd^*+ should create the following tree:



- (b) Implement your algorithm.
- 5.2. (a) Write an augmented, attributed grammar for a bottom-up parser that builds the syntax tree required for the previous exercise.
 (b) Rewrite your grammar to support infix expressions as input rather than postfix. Parenthesized subexpressions must be supported.
 (c) Implement both of the previous grammars using `occs`.
- 5.3. Using the LR parse table in Table 5.11 on page 369, show the contents of the state and symbol stacks as the following input is processed:

```
1 + ( 2 * 3 ) * 4
```

- 5.4. Create bottom-up, LALR(1) parse tables for the following grammar:

1.	<i>stmt</i>	\rightarrow	<i>expr SEMI</i>
2.			<i>WHILE LP expr RP stmt</i>
3.			<i>LC stmt_list RC</i>
4.	<i>stmt_list</i>	\rightarrow	<i>stmt_list stmt</i>
5.			ϵ
6.	<i>expr</i>	\rightarrow	<i>expr PLUS term</i>
7.			<i>STAR expr</i>
8.			<i>term</i>
9.			ϵ
10.	<i>term</i>	\rightarrow	<i>IDENTIFIER</i>
11.			<i>LP expr RP</i>

This grammar produces one shift/reduce conflict, which should be resolved according to the following table:

Input symbol	Token	Precedence	Associativity
()	LP RP	high	none
*	STAR	medium	right to left
+	PLUS	low	left to right
{ }	LC RC	none	left to right
;	SEMI	none	none

Hand in the tables and graphs showing both the LR and LALR state machines.

- 5.5. Modify the grammar in the previous exercise to include an assignment operator ($x=y;$) and an exponentiation operator ($x^{^y}$) is x to the y th power. Assignment should be the lowest precedence operator and exponentiation the highest. Make LALR(1) parse tables for this modified grammar.
- 5.6. Implement a parser that uses the parse tables created in Exercises 5.4 or 5.3.
- 5.7. Add code-generation actions to the grammar created in Exercises 5.4 or 5.3. You may output either assembly language or C.
- 5.8. Add the code-generation actions from the previous exercise to the parser implementation created in Exercise 5.5.
- 5.9. Why are fewer shift/reduce conflicts generated when imbedded actions are inserted into a production immediately to the right of terminal symbols?
- 5.10. Think up a real-world situation where a reduce/reduce conflict is inevitable (other than the one presented earlier).
- 5.11. A more-efficient method for creating LALR(1) parse tables than the one presented in the current chapter is described in [Aho], pp. 240–244. An LR(0) state machine is created and then lookaheads are added to it directly. You can see the basis of the method by looking at Figure 5.8 on page 366. Lookaheads appear on the tree in one of two ways, either they propagate down the tree from a parent to a child or they appear spontaneously. A lookahead character propagates from the item $[s \rightarrow \alpha . x \beta, c]$ to its child when $\text{FIRST}(\beta)$ is nullable. It appears spontaneously when $\text{FIRST}(!b)$ contains only terminal symbols. Modify `occs` to use this method rather than the one presented in the current chapter. Is the method in [Aho] actually more efficient once implementation details are considered? Why or why not?

- 5.12. An even more efficient method for creating LALR(1) parse tables, developed by DeRemer and Pennello, is described in [DeRemer79] and [DeRemer82]. It is summarized in [Tremblay], pp. 375–383. Modify **occ**s to use this table-creation method.
- 5.13. **LLama** and **occ**s currently supports optional and repeating productions, indicated with brackets:

```
s : a [optional] b  
| c [repeating zero or more times]* d  
;
```

These are implemented in the `start_opt()` and `end_opt()` subroutine presented in Chapter 4 and the transformations involved are discussed in Appendix E. **Occs** uses left-recursive lists and **LLama** uses a right-recursive lists for these transformations.

- (a) Add a + operator to **occ**s and **LLama** that can be used in place of the star in a repeating production. It causes the enclosed sub-production to repeat one or more times.
- (b) Add the operator `[...](a,b)*` for a to b repetitions of the sub-production— a and b are decimal numbers.
- (c) Add the `[...]*>` and `[...]+>` operators to **occ**s. These should work just like `[...]*` and `[...]+`, but should use right-recursive rather than left-recursive list productions.

Code Generation

Having looked at how parsers are created, it's now time to look at the real heart of the compiler: the code generation phase. We've actually been looking at code generation in a backhanded way. `LEX`, `LLama`, and `occs` are all compilers, after all, and they do generate code of sorts—the tables and drivers. This chapter looks at more conventional code generation and at the run-time environment by building a C compiler. I'm assuming that you can generalize from a description of a specific implementation once you understand the design issues. The grammar used here is summarized in Appendix C and is discussed piecemeal in the current chapter as it's used.

Because a production-quality, full ANSI compiler would both take up too much space and actually obscure some code-generation issues, I've implemented a proper subset of ANSI C. Most of the hard stuff is implemented, but I've omitted some functionality that you can add on your own if you're interested.¹ The modifications are left as exercises at the end of the chapter. Note that these limitations exist only at the semantic level. The grammar recognizes ANSI C, but the compiler ignores a few ANSI constructs when they are encountered. The compiler has the following limitations:

- Floating point is not supported.
- The `auto`, `const`, `volatile`, and `register` keywords are ignored.
- Structures do not form lvalues. You can't pass or return them by value or assign to a structure.
- Compile-time initializations work only on nonaggregate global variables like this:

```
int x = 5;
```

Arrays, structures, and local variables of any sort may not be initialized at compile time.

- Strong type checking on enumerated types is not supported. An enumerated type is treated just like an `int`, and any integer value can be assigned to a variable of that type, even if the value was declared as part of a different enumerated type. The

ANSI-C subset implemented.

1. Adding some of the missing stuff, like structure lvalues, floating point, and proper initializations, is a nontrivial enterprise. Other omissions, like bit fields, are easy.

elements of the enumeration list are treated just like a macro would be treated, it is just replaced by a `int` constant wherever it's found. The following declaration:

```
enum tag { Beaver, Wally, June, Ward; } Cleaver;
```

is treated like this:

```
#define Beaver 0
#define Wally 1
#define June 2
#define Ward 3

int Cleaver;
```

- Bit fields are ignored. The associated structure field is treated as if the bit field wasn't present in the input.
- Function prototypes are treated as simple `extern` declarations. The argument or type list is ignored. The new C++-style function-definition syntax is supported, however.
- One nonstandard escape sequence is supported in character constants: `\^C` (backslash-caret-letter) evaluates to the associated control character. It cannot be used in a string constant.

A simple, one-pass compiler is developed in this chapter: the compiler generates assembly language directly rather than creating a true intermediate language that is processed by a back end. I've taken this approach because it represents something of a worst-case scenario, and so makes a better example. The chapter starts out discussing the run-time environment and how a typical machine is used by the generated code. Intermediate languages are discussed briefly, and an output assembly language with a C-like syntax is developed. Internal data structures such as the symbol table are discussed, and then the compiler is developed.

One failing of the current chapter is that the recursive nature of the grammar precludes a strictly hierarchical discussion of the code-generation actions. As a consequence, there's a certain amount of skipping around in the grammar in the later sections. I've tried to minimize this as much as possible, but you may have to read the chapter more than once before the overall structure of the compiler becomes apparent. The entire grammar is listed hierarchically in Appendix C, and it will be worthwhile for you to study the grammar thoroughly before proceeding to the discussion of the code-generation issues.

As a final disclaimer, compilers are complicated programs and are notoriously difficult to debug. I don't expect the current compiler to be any different in this respect. The compiler has been tested as well as I am able, but there's nothing like having several thousand people reading your code to show up bugs that you never dreamed existed. The electronically distributed version will always be the most up-to-date, and bug fixes will be posted to USENET on a regular basis. Please notify me when you find a bug so I can tell everyone else about it (either c/o Software Engineering, or electronically; the addresses are in the Preface).

You must be thoroughly familiar with the bottom-up parse process described in the last chapter before continuing.

6.1 Intermediate Languages

Compilers often generate *intermediate languages* rather than translate an input file directly to binary. You can look at an intermediate language as a model assembly language, optimized for a nonexistent, but ideal, computer called a *virtual machine*. The

compiler's output can be tested by simulating this virtual machine with a computer program. There are several advantages to an intermediate-language approach. First, you can design an intermediate language with ease of optimization in mind, and thereby improve the resultant binary image because it can be more heavily optimized. Next, the intermediate-to-binary translation is usually done by a separate compilation pass called a *back end*, and you can provide several back ends for different target machines, all of which use the same parser and code generator (called the *front end*). By the same token, you can provide several front ends for different languages, all of which generate the same intermediate language and all of which share the same back end to generate real code.²

Back ends and front ends.

Virtual machines typically have many registers, very orthogonal instruction sets (all instructions can be performed on all registers, the syntax for register operations is the same as the memory-access syntax), and so forth. There are trade-offs, though. The ideal machine hardly ever maps to a real machine in an efficient way, so the generated code is typically larger and less efficient than it would be if the parser knew about the actual target machine. The ability to mix and match front and back ends compensates to some extent for the inefficiencies.

Virtual machine characteristics.

Intermediate languages typically take one of three forms. These are *triples*, *quads* (short for quadruples), and *postfix* (reverse-Polish) notation.

Most real assembly languages are made up of *triples* which are made of of three parts: an *operator*, a *source*, and a *destination* or *target*. For example, the 68000 instruction ADD.W D0,D1 is a triple that adds the contents of the D0 and D1 registers and puts the result in D1. A C representation of a similar triple would be:

```
d += s;
```

Triples (two-address instructions).

and a typical mathematical representation is:

$$(+=, d, s)$$

Triples are sometimes called *triplets* or *3-tuples*. They are also sometimes called *two-address* instructions because the binary representation of most instructions comprise an operator and source and destination addresses.

Quads, also called *quadruples* or *three-address* instructions, have four parts. The following quad has two sources, an operand, and a destination:

```
d = s1 + s2;
```

Quads (three-address instructions).

A more mathematical representation of the same quad would be:

$$(+, d, s1, s2)$$

Note that, in spite of the name, some quads, such as an assignment quad, have only three parts, with an empty fourth field. This empty element is usually marked with a dash:

$$(=, d, s, -)$$

The first field defines the operation. Since it's never empty, a dash there signifies subtraction.

Not all quads and triples involve implicit assignments. The first of the following triples compares two numbers and remembers the result of the comparison internally. The

Quads and triples without explicit assignments.

2. For example, the Microsoft Pascal, FORTRAN, and C compilers all share the same back end.

second triple, which branches to a specific label, is executed only if the previous comparison was true. Neither instruction involves an assignment.

```
(LESS_THAN, a, b)
(GOTO, target, -)
```

Arithmetic operations may not involve assignment either. The following two triples execute $A=B+C$. The first one does the addition and stores the result in an internal register called the *accumulator*. The second assigns the result of the previous operation—the $.-1$ —to A .

```
(+, B, C)
(=, A, .-1)
```

The dot represents the position of the current instruction: $.-1$ references the previous instruction, $.-2$ the instruction before that, and so on.

Advantages of quads versus triples.

Triples have one advantage over quads: they're very close in structure to many real assembly languages, so they can ease code generation for these assemblers. I'm using them in the compiler generated in this chapter for this reason. Quads have two advantages over triplets. They tend to be more compact; a quad like $(+, d, s1, s2)$ requires two triplets to do the same work:

```
(= d, s1);
(+= d, s2);
```

Also, certain optimizations are easier to do on quads because triples are more position dependent. In the previous example, the two triples must be treated as a single unit by any optimizations that move the code around, and the ordering of the two triples is important. Since it's self-contained, the quad can be moved more easily.

Postfix, Reverse-Polish Notation (RPN).

The third kind of common intermediate language is *postfix* or *Reverse-Polish Notation* (RPN.) Forth, PostScript, and Hewlett-Packard calculators are examples of postfix languages. This representation has several advantages. The first is that expressions can be evaluated with less work than usual: no parentheses are needed to represent them, and the compiler doesn't need to allocate temporary variables to evaluate postfix expressions—it uses a run-time stack instead. All operands are pushed onto the stack and all operators affect the top few stack elements. For example, an expression like this:

```
(1+2) * (3+4)
```

is represented in postfix as follows:

```
1 2 + 3 4 + *
```

and is evaluated as shown in Table 6.1.

Postfix languages often have various operators designed specifically for stack operations, such as a *dup* operator that duplicates and pushes the top of stack element— X^2 can be represented with:

```
X dup *
```

Another common stack operator is the *swap* operator, which swaps the two elements at the top of stack.

RPN has one other advantage: It's easy to reconstruct the original syntax tree from a postfix representation of the input language. (We'll look at this process further in Chapter Seven.) Since some optimizations require the production of a syntax tree, a postfix intermediate language is a convenient way to transfer a compiled program to the optimizer.

Table 6.1. Postfix Evaluation of $1 \ 2 \ + \ 3 \ 4 \ + \ *$

stack	input	comments
<i>empty</i>	$1 \ 2 \ + \ 3 \ 4 \ + \ *$	<i>push 1</i>
1	$2 \ + \ 3 \ 4 \ + \ *$	<i>push 2</i>
1 2	$+ \ 3 \ 4 \ + \ *$	<i>add the two items at top of stack and replace them with the result.</i>
3	$3 \ 4 \ + \ *$	<i>push 3</i>
3 3	$4 \ + \ *$	<i>push 4</i>
3 3 4	$+ \ *$	<i>add the two items at top of stack and replace them with the result.</i>
3 7	*	<i>multiply the two items at top of stack and replace them with the result.</i>
21		<i>the result is at top of stack.</i>

6.2 C-code: An Intermediate Language and Virtual Machine³

The syntax of an intermediate language is pretty much arbitrary. Ideally it reflects the characteristics of the most likely target machines. For purposes of clarity, however, I've decided to use a C-like intermediate language in the compiler created later in this chapter. I've defined a C subset in which all instructions have direct analogs in most assembly languages—all the instructions translate directly into a small number of machine instructions (usually one). This way, the assembly-language syntax will be familiar to you, regardless of your background. In addition, you can use your C compiler and normal debugging tools to exercise the output from the compiler. The code is simple enough so that a translation to assembly language is very straightforward. I've dubbed this intermediate language *C-code*.

C-code is really more of an assembly language than a true intermediate language. The main problem is that it is very restrictive about things like word width, alignment, and storage classes. It forces the compiler itself to worry about details that are usually handled by the back end and, as a consequence, makes a back end harder to write because the back end must occasionally undo some of the things that the compiler has done (like add padding to local variables to get proper alignment).

The current section describes C-code in considerable depth. Though the description is, by necessity, full of specific implementation details, reading it should give you a good idea of the sorts of things that are required in a more general-purpose intermediate language. The current section also serves as a review of assembly-language concepts. The C-code description is pretty terse in places, and a previous familiarity with a real assembly language will help considerably in the following discussion. Finally, I've also used the current section as a vehicle for discussing the memory organization and subroutine-linkage procedures that are almost universal among C compilers, so you should read through it even if you're very familiar with assembler.

The C-code virtual machine—the hypothetical machine that would run C-code as its assembly language—is modeled by a series of macros in a file called `<tools/virtual.h>`,

`<tools/virtual.h>`

3. For convenience, all the C-code directives described in this section are also summarized in Appendix F.

which should be **#included** at the top of every C-code file. This file contains definitions that allow you to compile and run the generated C-code, using your normal C compiler as an assembler. *Virtual.h* is described piecemeal as the various C-code language elements are discussed. Unfortunately, the code in *virtual.h* is not particularly portable. I've made all sorts of assumptions about the sizes of various data types, the ordering of bytes within a word, and so forth. You can compensate for these system-dependent problems by modifying the include file, however.

A typical development cycle looks like this:

<code>vi file.c</code>	<i>Edit the C source-code file.</i>
<code>c file.c</code>	<i>Run the source code through the compiler developed in the current chapter. The compiler generates a C-code output file called <code>output.c</code>.</i>
<code>cc output.c</code>	<i>"Assemble" the compiler's output, using your normal compiler as an assembler. You can debug the compiler's output using dbx, CodeView, or whatever debugger you normally use.</i>

Table 6.2 shows how a C input file is translated into C-code by the compiler in the current chapter. The rightmost column demonstrates how the C-code relates to 8086 assembler. Table 6.2 is intended to demonstrate what's happening in a general way. Don't worry about the details—I'll discuss what's going on in great detail as the chapter progresses. For now, notice how similar the C-code is to the true assembly language. It's easy to translate from one to another. C-code is really an assembly language, regardless of the superficial similarities to C.

6.2.1 Names and White Space

White space.

As in C, white space is ignored in C-code except as needed to separate tokens. Comments delimited with `/* . . . */` are treated as white space. Multiple-line comments are not permitted—the `/*` and `*/` must be on the same line.

Identifiers.

Identifiers can be made up of letters, digits, and underscores only. The first character in the name may not be a digit, and names are restricted to 31 characters. Names are truncated if they're longer. In addition to the standard C keywords, all of the C-code directives discussed below should be treated as keywords and may not be used for identifiers.

6.2.2 Basic Types

byte, word, lword, ptr, array, record.

The primary design considerations in the virtual machine are controlled by the sizes of the variables that are manipulated by the machine. Four basic data types are supported: 8-bit bytes, 16-bit words, 32-bit long words, and generic pointers (nominally 32 bits). These types are declared with the `byte`, `word`, `lword`, and `ptr` keywords respectively. The `array` and `record` keywords can be used as a synonym for `byte` in order to provide some self-documentation when declaring a structure or array of structures. All but `ptrs` are signed quantities.

`<tools/c-code.h>`

Listing 6.1 shows the type definitions from *virtual.h*. I'm making nonportable assumptions about word widths here; you may have to change these definitions in your own system. The `ptr` type is a character pointer, as compared to a `void` pointer, so that it can be incremented if necessary.

The `<tools/c-code.h>` file, which is **#included** on line one of Listing 6.1, contains various definitions that control the widths of basic types. It is shown in Listing 6.2. This information has been split into a second header file because it is likely to be used by the compiler that is generating C-code, but is not much use in interpreting the C-code.

Table 6.2. Translating C to C-code to 8086 Assembly Language

C Input	C-Code Output	8086 Assembler
	#include <tools/virtual.h> #define T(x) SEG(bss) #define L0 1 /* strcpy: locals */ #define L1 2 /* strcpy: temps. */ #undef T #define T(n) (fp-L0-(n*4)) SEG(code)	_BSS SEGMENT WORD PUBLIC
strcpy(dst, src)	PROC(_strcpy,public)	_BSS ENDS
char *dst, *src;	/* fp+4 = dst [argument] */ /* fp+8 = src [argument] */ link(L0+L1);	_TEXT SEGMENT WORD PUBLIC ASSUME CS:_TEXT
{		PUBLIC _strcpy _strcpy PROC NEAR
char *start;	/* fp-4 = start [variable] */	; [bp+2] = dst ; [bp+4] = src push bp mov bp,sp sub sp,6 ; [bp-2] = start
start = dst;	BP(fp-4) = BP(fp+4);	mov ax,WORD PTR [bp+2] mov WORD PTR [bp-2],ax
while(*src)	TST1: EQ(*BP(fp+8),0) /* src */ goto EXIT1;	TST1: mov bx,WORD PTR [bp+4] mov al,BYTE PTR [bx] or al,al jz EXIT1
*dst++ = *src++ ;	BP(T(1)) = BP(fp+4); /* t1=dst */ BP(fp+4) += 1; /* dst++ */ BP(T(2)) = BP(fp+8); /* t2=src */ BP(fp+8) += 1; /* src++ */ *BP(T(1)) = *BP(T(2)); /* *t1=t2 */	mov si,WORD PTR [bp+2] inc WORD PTR [bp+2] mov di,WORD PTR [bp+4] inc WORD PTR [bp+4] mov bx,di mov al,BYTE PTR [bx] mov bx,si mov BYTE PTR [bx],al jmp TST1
}	goto TST1; EXIT1:	EXIT1: mov bx,WORD PTR [bp-4] jmp RET1
return start;	rF.pp = BP(fp-4); goto RET1;	
}	RET1: unlink(); ret(); ENDP(_strcpy)	RET1: mov sp,bp _strcpy ENDP _TEXT ENDS END

6.2.3 The Virtual Machine: Registers, Stack, and Memory

The C-code virtual machine is pictured in Figure 6.1. The machine has a set of 16 general-purpose *registers* named *r0*, *r1*, *r2*, and so forth. An entire 32-bit register can be accessed from any instruction, as can either of the 16-bit words or any of four 8-bit bytes that comprise the same register. These registers are memory locations that are physically part of the CPU itself—they don’t have addresses. Use the syntax shown in Table 6.3 to access a register. The name must always be fully qualified; one of the forms shown in Table 6.3 must always be used. The register name by itself (without the dot and type reference) is illegal.

The register set.

Listing 6.1. *virtual.h*— Basic Types

```

1 #include <tools/c-code.h>
2                                         /* Basic types           */
3 typedef char   byte;                  /* 8 bit                */
4 typedef short  word;                 /* 16 bit               */
5 typedef long   lword;                /* 32 bit               */
6 typedef char   *ptr;                 /* Nominally 32 bit.    */
7                                         /* Aliases for "byte." */
8 typedef byte   array;                /*                         */
9 typedef byte   record;

```

Listing 6.2. *c-code.h*— Various Widths

```

1 #define BYTE_WIDTH      1      /* Widths of the basic types.          */
2 #define WORD_WIDTH      2
3 #define LWORD_WIDTH     4
4 #define PTR_WIDTH       4
5
6 #define BYTE_HIGH_BIT   "0xff80"      /* High-bit mask. */
7 #define WORD_HIGH_BIT   "0x8000"
8 #define LWORD_HIGH_BIT  "0x80000000L"

```

The stack, fp, sp.

In addition to the register set, there is a 1024-element, 32-bit wide stack, and two special-purpose registers that point into the stack: the *fp* and *sp* registers—discussed in depth, below.

The instruction pointer,
ip.

Finally, there is the *ip* or instruction-pointer register. This register holds the address (in the code segment) of the next instruction to execute, not of the current instruction. It is updated every time an instruction is processed, and is modified indirectly by various instructions. A *call*, for example, pushes the *ip* and then transfers control to somewhere else. A *ret* pops the address at top of stack into the *ip*. A *goto* modifies the *ip* directly. The instruction pointer is not used directly by the current compiler, and its value won't change if you just use your C compiler to assemble the C-code output. Access to it is occasionally useful, however, and the register is available in all real machines. It's included here for completeness' sake.

Virtual machine imple-
mentation.

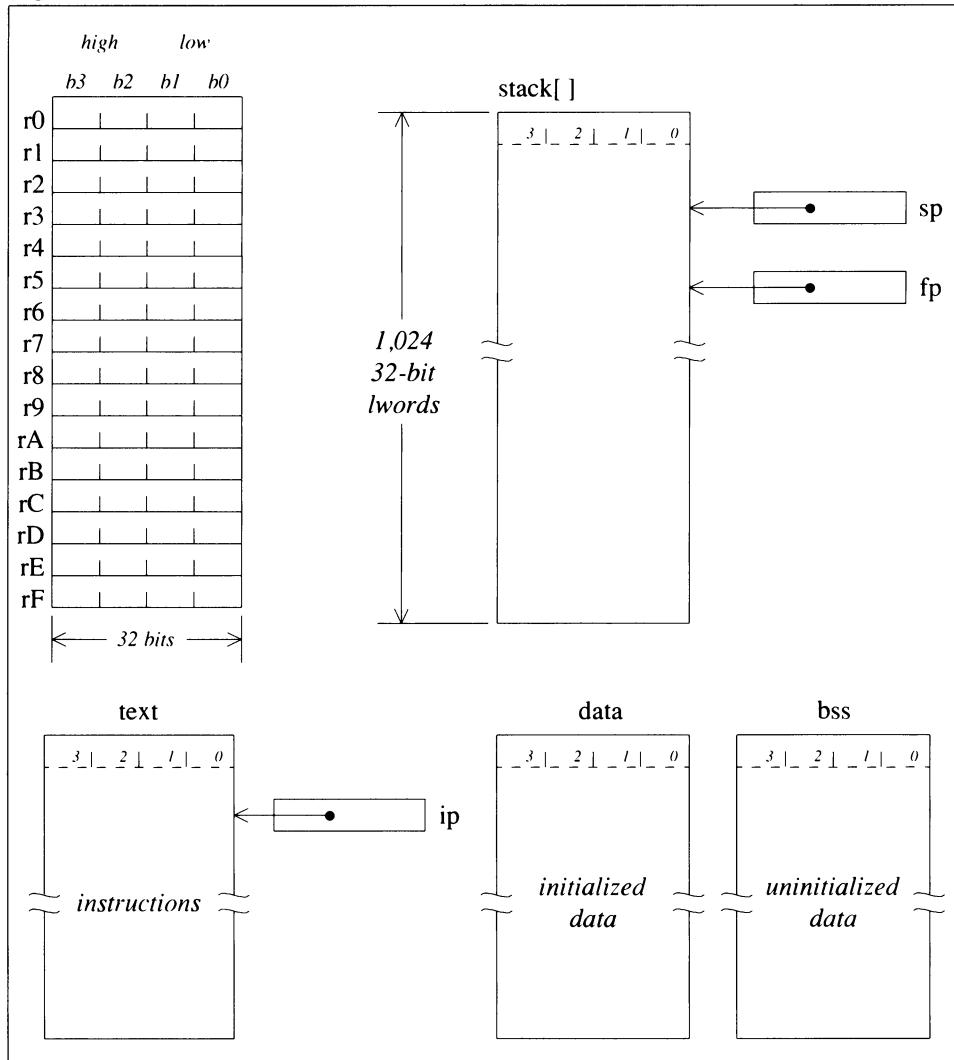
The register set and stack are implemented in Listings 6.3 and 6.4. Note that the definition for *reg* on lines 21 to 28 of Listing 6.4 is not portable because I'm assuming that an 8086-style byte ordering is used—the least-significant byte is lowest in memory, and the most-significant byte is highest. Exactly the opposite holds in 68000-style machines. The LSB is highest in memory, so you'd have to redefine the *_words* and *_bytes* fields as follows for those machines:

```

struct _words { word high, low; };
struct _bytes { byte b3, b2, b1, b0; };

```

The address of an object is the physical address of the least-significant byte in both the 8086 and 68000 family. Other machines may require even more shuffling.

Figure 6.1. The C-code Virtual Machine**Table 6.3.** The Virtual-Machine Register Set

rl contains:	access syntax:
pointer	rl.pp
32-bit long word	rl.l
two 16-bit words	rl.w.high rl.w.low
four 8-bit bytes (byte 0 is low)	rl.b.b3 rl.b.b2 rl.b.b1 rl.b.b0

Listing 6.3. *c-code.h*— Stack-size Definitions

```

9 #define SWIDTH          LWORD_WIDTH      /* Stack width (in bytes).      */
10 #define SDEPTH           1024            /* Number of elements in stack. */

```

Listing 6.4. *virtual.h*— The Register Set and Stack

```

10 #ifdef ALLOC
11 #   define I(x)  x
12 #   define CLASS /* empty */
13 #else
14 #   define I(x) /* empty */
15 #   define CLASS extern
16 #endif
17
18 struct _words { word  low, high; };
19 struct _bytes { byte  b0, b1, b2, b3; };      /* b0 is LSB, b3 is MSB */
20
21 typedef union reg
22 {
23     char          *pp;             /* pointer */
24     lword         l;              /* long word */
25     struct _words w;             /* two 16-bit words */
26     struct _bytes b;             /* four 8-bit bytes */
27 }
28 reg;
29
30 CLASS reg r0, r1, r2, r3, r4, r5, r6, r7;      /* Registers */
31 CLASS reg r8, r9, rA, rB, rC, rD, rE, rF;
32
33 CLASS reg      stack[ SDEPTH ];                  /* run-time stack */
34 CLASS reg      *__sp I(= &stack[ SDEPTH ]);    /* Stack pointer */
35 CLASS reg      *__fp I(= &stack[ SDEPTH ]);    /* Frame pointer */
36
37 #define fp      ((char *) __fp)
38 #define sp      ((char *) __sp)

```

Note that the stack itself is declared as an array of `reg` unions on line 33 of Listing 6.4, and the stack and frame pointers are `reg` pointers. The `sp` and `fp` registers are referenced in C-code directives using the macros on lines 37 and 38, however. The cast assures that pointer arithmetic is defeated in the based addressing modes discussed below.

The stack is deliberately made as wide as the worst-case basic type—if a 64-bit `double` were supported, I would have made the stack 64 bits wide. This way, all stack access can be done with a single instruction. I've done this in order to make the back end's life a little easier. It's a simple matter to translate single push and pop instructions into multiple instructions (if the target-machines stack is 16 bits wide, for example). It's difficult to go in the other direction, however—to translate multiple pushes and pops into single instructions.

You must put the following two lines into only one module of your program to actually allocate space for the stack and register set:

```

#define ALLOC
#include <tools/virtual.h>

```

Allocating space for the virtual-register set, ALLOC.

You can create a two-line file, compile it, and link it to your other code if you like.

When ALLOC is defined, the definitions on lines 11 and 12 are activated. Here, CLASS evaluates to an empty string, so all invocations of the CLASS macro effectively disappear from the file. The I () macro evaluates to its argument, so an invocation like

CLASS.

The I() macro.

```
I(= &stack[SDEPTH]);
```

evaluates to

```
= &stack[SDEPTH];
```

The macro's argument is the entire string =&stack[SDEPTH]. When ALLOC is not defined, the opposite situation holds, CLASS expands to the keyword **extern** and the contents of the I () macro are discarded.

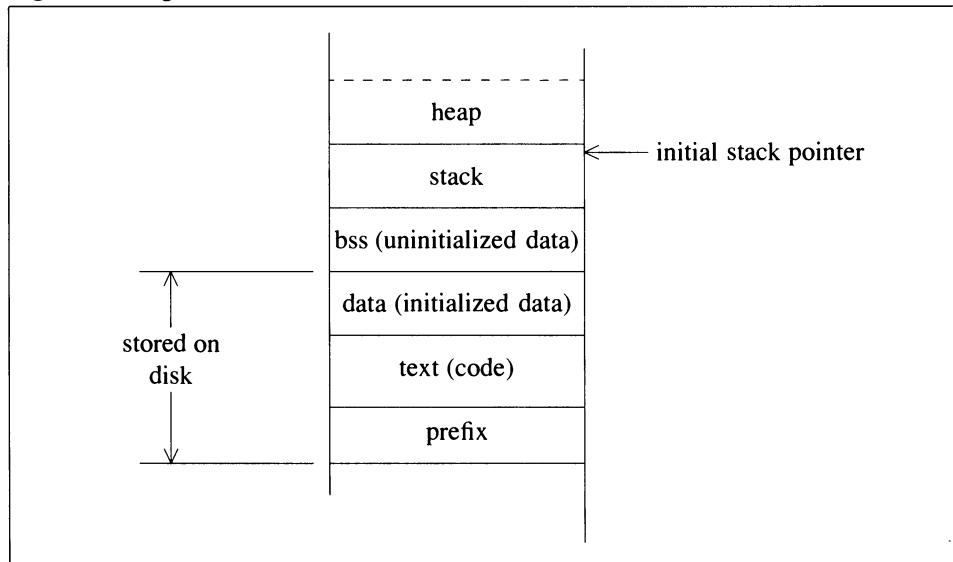
The run-time stack, stack pointer, and frame pointer registers are defined at the bottom of Listing 6.4. The stack is the same width as a register, and the stack pointer is initialized to point just above the top stack element. It grows towards low memory with the first push.

6.2.4 Memory Organization: Segments

The stack represents only a portion of the memory that can be used by the running program. This memory is partitioned into several *segments*, one of which is the *stack segment*. Figure 6.2 shows how the segments are usually arranged.

Segments.

Figure 6.2. Segments



The rationale behind this partitioning is best understood by looking at how a program is loaded by the operating system. Initially, a contiguous region of memory is allocated for the program to use. This region can be physical memory, as it is in most microcomputers, or can be virtual memory, as on a mainframe. In the later case, the physical memory might not form a contiguous block, but it's convenient to look at it that way.

The program-load process.

The *prefix* segment holds information that is used by the operating system to load the program. Typically, the sizes of the various segments are stored there, as are things like the initial values of the program counter and stack pointer. Some operating systems (like MS-DOS) read the prefix into memory, and it's available for the program to use at run time. Other operating systems read the prefix to get the information that it contains, and then overwrite the prefix with the other segments.

The program prefix.

Executable image.

Text (or code) and (initialized) *data* segments.

Bss segment.

Stack segment.

Heap segment.

Dynamic arrays.

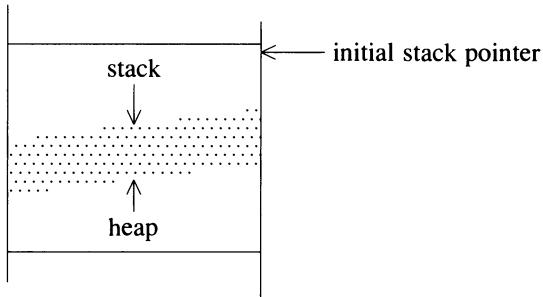
Combined heap and stack.

Using the information in the prefix, the operating system then allocates a block of memory large enough to hold the remainder of the program, called the *executable image*. The *text* and *data* segments are then copied directly from the disk into memory. The *text* or *code* segment holds the executable code; the *data* segment holds initialized data, only. That's why an initialized **static** local variable comes up with an initial value, but once that value is changed it stays changed—the initial value is just read from the disk into the proper place in memory. Note that there are two types of initialized data: variables that have been given an initial value when declared, and initialized constants (like string constants) whose values are not expected to change. The constant data is sometimes stored in the *text* segment, along with the code, in order to make the operating system's life a little easier. It knows that the *text* segment is not modified, so it doesn't have to swap this region out to the disk if it needs the memory for another process.

One of two things can happen once the *text* and *data* segments are loaded, depending on the operating system. Either control can be transferred directly to the program, which must continue the initialization process, or the operating system itself initializes the other segments. In any event, the *bss*⁴ segment, which holds all uninitialized data, is initialized to all zeros at load-time. (There's no point in storing masses of zeros on the disk.) The stack pointer is then initialized to point at the correct part of the *stack* segment, which holds the run-time stack used for subroutine calls, and so forth. Stacks grow down in most computers, so the stack pointer is typically initialized to the top of the stack segment. Various pointers that manipulate the *heap*, the region of memory used for dynamic storage, are also initialized. The heap is used, in C, by the memory-allocation subroutines `malloc()` and `free()`.⁵ In some languages the heap is also used in code generated by the compiler itself. The C++ new operator allocates memory for an object; and a user-defined class might do memory allocation transparently when an object in that class is declared. Similarly, PL/1 supports dynamically allocable arrays—arrays whose size is not known until run time, and the PL/1 compiler uses the heap to allocate space for that array when the scope of that declaration is entered at run time (when the block that holds the declaration is entered). The compiler just translates a compile-time declaration into code that allocates space for the array and initializes a pointer to access the first element—it effectively calls `malloc()` when the subroutine is entered and `free()` at exit. The contents of the stack and heap segments are typically uninitialized at load time, so the contents of variables allocated from these regions are undefined. The heap is typically at the top of the image because it may have to grow larger as the program runs, in which case the memory-allocation functions request that the operating system enlarge the size of the executable image. The heap and stack segments are often combined, however. Since the stack pointer typically grows down, and the memory allocation from the heap can start in low memory and go up, the shared space can be allocated from both ends as needed, like this:

4. *Bss* stands for “block starting with symbol,” a term dating from the Mesozoic, at least.

5. The source code for `malloc()` and `free()` implementation is in [K&R] pp. 185–189.



This way, if a program has an unusually large stack and unusually small heap or vice versa, space can be transferred from one area to the other with ease. It's difficult to expand the size of the heap after the top of heap runs into the stack, however, and such expansion almost always results in a fragmented heap. The virtual-memory manager could solve the fragmentation problem, but it might not.

Many compilers provide a mechanism for the program to determine the positions of various segments at run time. For example, the UNIX compiler automatically declares a variable at the end of the *bss* segment called *ebss*. All memory whose address is greater than *&ebss* must be in the stack or heap segments. Other variables (called *etext* and *edata*) are provided for the other segment boundaries. Similarly, many compilers generate code that checks the value of the stack pointer when every subroutine is entered and terminates the program if the stack pointer is not in the stack segment.

C-code lets you change from one segment to another by issuing a `SEG()` directive, which takes one of the following forms:

```
SEG( text )
SEG( data )
SEG( bss )
```

Note that no semicolon is used here—this is the case for all C-code directives whose names are in all caps. Everything that follows a `SEG(text)` (up to the next `SEG()` directive or end of file) is in the *text* segment, and so forth. You can use `SEG(code)` instead of `SEG(text)`. There's no direct access to the *prefix*, *stack*, and *heap* segments because these exist only at run time, and the compiler uses other mechanisms to access them. You may not change segments in the middle of a subroutine. No `SEG` directives can appear between the `PROC` and `ENDP` directives, discussed below. The `SEG()` directive is defined as an empty macro in *virtual.h*. It's shown in Listing 6.5.

Segment-end markers:
ebss, *etext*, *edata*.

`SEG()` directive.

No semicolon used with upper-case C-code directives.

Listing 6.5. *virtual.h*—The `SEG()` directive.

```
39 #define SEG(segment) /* empty */
```

6.2.5 Variable Declarations: Storage Classes and Alignment

C-code supports global-level variable declarations only. All variables must be declared in the *data* or *bss* segments, depending on whether or not they are initialized. Four storage classes are available (they are related to C in Table 6.4):

`private` Space is allocated for the variable, but the variable cannot be accessed from outside the current file. In C, this class is used for all `static` variables, be they local or global. Initialized variables go into the *data* segment, other variables go into the *bss* segment.

public	Space is allocated for the variable, and the variable can be accessed from any file in the current program. In C, this class is used for all initialized nonstatic global variables. It is illegal for two public variables in the same program to have the same name, even if they're declared in different files. Since public variables must be initialized when declared, they must be in the <i>data</i> segment.
common	Space for this variable is allocated by the linker. If a variable with a given name is declared common in one module and public in another, then the public definition takes precedence. If there are nothing but common definitions for a variable, then the linker allocates space for that variable in the <i>bss</i> segment. C uses this storage class for all uninitialized global variables.
external	Space for this variable is allocated elsewhere. If a label is external, an identical label must be declared common or public in some other module of the program. This storage class is not used for variables in the current application, all of which are common, public, or private. It is used for subroutines, though.

Table 6.4. Converting C Storage Classes to C-code Storage Classes

	static	not static	
		definition	declaration (extern or prototype)
Subroutine (in <i>text</i> segment).	private	public	external
Uninitialized variable (in <i>bss</i> segment).	private		common
Initialized variable (in <i>data</i> segment).	private		public

The following example shows how these classes are used in C:

```
int kookla = 1;      /* Public--it's initialized and not static. */
static int fran;    /* Private--it's static. Initialized to 0. */
int ollie;          /* Common--it's declared without an extern. */
int ollie;          /* Not a redefinition because default class */
                    /* is extern. Both instances of ollie are */
                    /* common. */
void ronnie()
{
    static float george; /* Private--it's declared static. */
extern dan;           /* Common. */
    int dick = 1;        /* No declaration is generated. (The */
                        /* memory is allocated on the stack). */
}
```

The four storage classes can be broken into two broad categories—classes that allocate space in the executable image (as stored on the disk) and classes that don't allocate space. I'll have to dip into real assembly language (8086 assembler) for a moment to show what's actually going on. The assembler actually builds a physical copy of the executable image in memory as it works, and it copies this image to the disk when assembly is complete. The assembler keeps track of the current position in the executable image that it's building with an internal variable called the *location counter*. A directive like the following allocates space for a variable:

```
_var : dw 10
```

This instruction tells the assembler to do the following:

- Create a symbol-table entry for `_var` and remember the current value of the location counter there.
- Fill the next two bytes with the number 10. Two bytes are allocated because of the `dw`—other codes are used for different sizes; `db` allocates one byte, for example. The 10 is taken from the `dw` directive.
- Increment the location counter by two bytes.

The allocated space (and the number 10, which is used here to fill that space) end up on the disk as part of the executable image. Instructions are processed in much the same way. For example, a

```
MOV ax, _var
```

instruction moves the contents of `_var` into the `ax` register. This instruction tells the assembler to do the following:

- Copy a binary code representing the “move into `ax`” operation into the place referenced by the current location counter. This binary code is called the *op code*.
op code.
- Copy the address of `_var` into the next two bytes. This address is the location-counter value that was remembered in the symbol table when space for `_var` was allocated.
- Increment the location counter by three.

From the assembler’s perspective, code and data are the same thing. All it knows about is the current location counter. There is nothing preventing us from putting a `MOV` instruction into memory as follows:

```
db 0A0H      ; 0xA0 is the 8086 op code for "MOV into AX"
dw _var       ; address of _var
```

Applying the foregoing to C-code storage classes, `public` and `private` definitions are translated into `dw` directives in 8086 assembler. They cause the location counter to be moved, and some value is copied into the allocated space. This value ends up on the disk as part of the executable image. The `external` directive is at the other extreme. No space is allocated; the location counter is not modified. The assembler does create a symbol-table entry for the associated label, however. If that label is referenced in the code, place holders are put into the image instead of the actual addresses, which are not known by the compiler. The linker replaces all place holders with the correct addresses when it puts together the final program. The `external` directive must reference a region of memory allocated with a `dw` directive or equivalent somewhere else in the program. The unpatched binary image (with the place holders still in it) is usually called a *relocatable object module*.⁶

Relocatable object module.

The `common` storage class is somewhere in between `external` and `public`. If the name associated with the `common` is used elsewhere in a `public` or `private` declaration, then the `common` is treated just like an `external`. Things change when all

6. One common way to organize a relocatable module puts a symbol table into the object file. The table has one element for each unresolved reference—`private` symbols are not put into the table. The symbol-table elements hold the symbol’s name and the offset from the start of the file to the first place holder that references that symbol. That place holder, in turn, holds the offset to the next place holder for the same symbol, in a manner similar to a linked list. The last element of the list is usually marked with zeros. Most 8086 relocatable-object-module formats derive from the specification described in [Intel] and in [Armbrust].

references to the variable are commons, however. Space is allocated for the variable at load time, not at compile time. Remember, the executable image on the disk represents only part of the space that is used when the program is running. Space for a common is always allocated in the *bss* segment, which is created when the program is loaded—it has no existence at compile time and won't be part of the binary image on the disk. The linker replaces the place holders in those instructions that use the common variable with references to the place at which the variable is found at run time, but no dw directive is generated. Space is allocated implicitly because the linker leaves a note in the executable image's file header that tells the loader the size of the *bss* region. Instead of incrementing the location counter at compile time, the assembler tells the loader to increment it at load time, by making the size of the executable image large enough to encompass the extra space.

Allocating variables.

Returning to C-code, all variables must be declared using a basic type (byte, word, lword, or ptr) and one of the foregoing storage classes (private, public, common, external). In addition, one-dimensional arrays can be declared using trailing brackets. Multi-dimensioned arrays are not permitted. The following declarations (only) are available:

```
class type name;           /* single variable */
class type name [ constant ]; /* array */
```

A structure must be declared as a byte array, though the keyword record can be used as a synonym for byte for documentation purposes.

Public and private variables in the *data* segment may be initialized with an explicit C-style initializer. Character constants and implicit array sizes are both supported. For example:

```
SEG( data )
public byte name      = 'z' ;
public byte name[]   = 'a', 'b', 'c', 'd', '\0' ;
public byte name[]   = "abcd";
private word name     = 10;
public word name[3]  = 10, 11, 12 ;
```

The double-quote syntax can be used only to initialize byte arrays. A C declaration such as the following:

```
kings[4] =
{
    "henry",
    "kong",
    "elvis"
    "balthazar"
};
```

must be declared as follows in C-code:

```
SEG( data );
private byte S1[]      = "henry\0"      ;
private byte S2[]      = "kong\0"       ;
private byte S3[]      = "elvis\0"       ;
private byte S4[]      = "balthazar\0"  ;
public  ptr  kings[4] = { S1, S2, S3, S4 };
```

These declarations are all in the *data* segment because they're initialized. The anonymous string names must be private because the same labels may be used for other anonymous strings in other files. kings is public, because it wasn't declared **static**. The *virtual.h* definitions for the various storage classes are in Listing 6.6.

Listing 6.6. *virtual.h*—Storage Classes

```

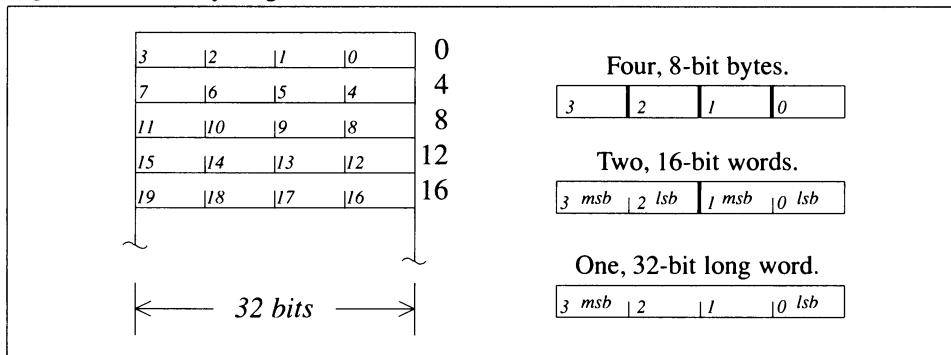
40 #define public           /* empty */
41 #define common          /* empty */
42 #define private         static
43 #define external        extern

```

The C-code assembler assumes that memory is made up of 32-bit wide cells, each of which occupies four addresses. The various basic types control the way that the four bytes can be accessed. In particular, the least significant byte of an object must also be at an address that is an even multiple of the object's size—a byte can be anywhere, a word must be at an even address, an lword and ptr must be at an address that's an even multiple of four. All objects are stored with the least-significant byte at the lowest physical address and the most-significant byte at the highest address. A pointer to an object holds the physical address of the least-significant byte. The system is pictured in Figure 6.3.

Alignment.

LSB is at the lowest address.

Figure 6.3. Memory Alignment

A good analogy for these restrictions is a book printed in a typeface that's half a page high. The typeface is also so wide that only two digits can fit on a line. The rules are that a two-digit number must be on a single line. A four-digit number, which requires two lines to print, must be on a single page. An eight-digit number, which requires four lines to print, must be on two facing pages.

These restrictions are called *alignment* restrictions. (You would say that words must be aligned at even addresses, and so forth.) Alignment restrictions are usually imposed by the hardware for efficiency's sake. If the data path from memory to the CPU is 32 bits wide, the machine would like to fetch all 32 bits at once, and this is difficult to do if the 32-bit word can be split across a multiple-of-four boundary.

The actual byte ordering is not critical in the present application provided that it's consistent, but the alignment restrictions are very important when doing something like allocating space for a structure. For example, the following structure requires 16 bytes rather than nine, because the 16-bit **int** must be aligned on an even boundary; the 32-bit **long** must start at an address that's an even multiple of four, and padding is required at the end to assure that the next object in memory is aligned properly, regardless of its type. A declaration like the following:

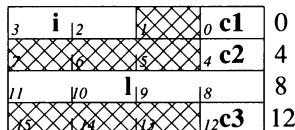
Structure-field allocation and alignment.

```

struct
{
    char   c1;
    int    i;
    char   c2;
    long   l;
    char   c3;
}
fred;

```

allocates fields placed in memory like this:



Though C guarantees that the ordering of structure members will be preserved, a clever compiler for another language might shuffle around the ordering of the fields to make the structure smaller. If `c1` were moved next to `c2`, then four bytes would be saved. The three bytes of padding at the end of the structure are required in case you have an array of structures. A clever compiler could actually eliminate the padding in the current example, but nothing could be done if the first field of the structure was a `long`, which needs to be aligned on a multiple-of-four boundary.

`ALIGN()` directive.

C-code variable declarations automatically force proper alignment for that variable, so if you declare a byte-sized variable followed by a 16-bit word, one byte of memory is wasted because the word requires even alignment. There are occasional situations where alignment must be forced explicitly, and C-code provides the `ALIGN(type)` directive for this purpose. The assembler inserts sufficient padding to assure that the next declared variable is aligned as if it were of the indicated *type*. `ALIGN` is defined in *virtual.h* in Listing 6.7. The worst-case alignment boundary is declared for use by the compiler in *c-code.h* (Listing 6.8).

Listing 6.7. *virtual.h*—The `ALIGN` Directive

```
44 #define ALIGN(type) /* empty */
```

Listing 6.8. *c-code.h*—Worst-case Alignment Restriction

```
11 #define ALIGN_WORST LWORD_WIDTH /* Long word is worst-case alignment. */
```

6.2.6 Addressing Modes

Constants and all variable and register names are referenced in instructions by means of various *addressing modes*, summarized in Table 6.5.

Immediate mode.

The *immediate* addressing mode is used for numbers. C-style hex, octal, decimal, and character constants are all recognized (0xabc 0377 123 'c'); lword (32-bit) constants should be indicated by a trailing `L` (0x12345678L). The immediate addressing mode is used in real assemblers to put a physical number into an instruction rather than an address. Many real assemblers require a special symbol like `#` to precede immediate data.

Direct mode.

The *direct* mode is used to fetch the contents of a variable. There is usually no special operator here; you just use the variable's name. Note, however, that array and

function names cannot be accessed in direct mode—their name is always treated as the address of the first element or instructions.

The *effective-address* mode gets the address of an object. The operand of the & operator is usually a label (&x where x is a name in a previous declaration), but it can also be used in conjunction with one of the based addressing modes, discussed shortly. Note that the & should be omitted from a function or array name.

The *indirect* modes work only on objects of pointer type: a variable that was declared `ptr`, a register name followed by the .pp selector, or the `sp` or `fp` registers. You must do two things to access an object indirectly. You must surround the variable or register that holds the address with parentheses to indicate indirection, and you must tell the C-code assembler the kind of object that the pointer is referencing. You do the latter by preceding the parenthesis with one of the following one- or two-character modifiers:

Indirect modes.

Indirection indicated by parentheses.

Code	Points at object of this type.
B	byte
W	word
L	lword
P	ptr
BP	pointer to byte
WP	pointer to word
LP	pointer to lword
PP	pointer to ptr

For example, `W(r0.pp)` fetches the word whose address is in `r0.pp`. You can also access an object at a specified offset from a pointer with the following syntax:

`W(p + offset)`

The *offset* may be a number, a numeric register reference (as compared to a pointer reference), or a numeric variable. The following example references the word whose address is derived by adding together the contents of the `fp` register and `r0.w.low`:

`W(fp + r0.w.low)`

The following instruction references the long word at an offset of -16 (bytes) from the frame pointer:

`L(fp-16)`

If the fetched object is of type BP, WP, LP, or PP, a star may be added to fetch the object to which the pointer points. For example, `WP(fp+6)` fetches the word pointer at offset 6 from the frame pointer; `*WP(fp+6)` fetches the object pointed to by that word pointer. Note that these *double indirect* modes, though found on most big machines like VAX's, are missing from many machines. You would have to use two instructions to access the referenced object—like this:

`WP(p), *WP(p)`

```
r0.pp = WP(fp+6);
r1.w  = W(r0.pp);
```

The effective-address and indirect modes can be combined in various ways, summarized in Table 6.6. The alternate syntaxes are useful when doing code generation, because variables on the stack can be treated the same way as variables at fixed addresses. For example, the address of a word pointer that is stored at offset 8 from the frame pointer can be fetched with `&WP(fp+6)`, and the address of a word pointer, `_p`, at a fixed address can be accessed with `&WP(&_p)`. The convenience will become evident when code-generation for the arithmetic operators is discussed, below. You can look at

Combined effective-address and indirect modes. `&WP(&p)`.

Table 6.5. Addressing Modes

Mode	Example	Notes
immediate	93	Decimal number. Use leading 0x for hex, 0 for octal.
direct	x, r0.1	Contents of variable or register.
indirect	B(p) W(p) L(p) P(p) BP(p) WP(p) LP(p) PP(p)	byte whose address is in p. word whose address is in p. lword whose address is in p. ptr whose address is in p. byte pointer whose address is in p. word pointer whose address is in p. lword pointer whose address is in p. ptr pointer whose address is in p.
double indirect	*BP(p) *WP(p) *LP(p) *PP(p)	byte pointed to by pointer whose address is in p. word pointed to by word pointer whose address is in p. lword pointed to by lword pointer whose address is in p. ptr pointed to by ptr whose address is in p.
based indirect	B(p+N) W(p+N) L(pN) ...	byte at byte offset N from address in p. word at byte offset N from address in p. lword at byte offset N from address in p.
effective address	&name &W(p+N)	Address of variable or first element of array. Address of word at offset +n from the pointer p (The effective-address modes can also be used with other indirect modes see below.)
A generic pointer, <i>p</i> , is a variable declared <i>ptr</i> or a pointer register: rN.pp, fp, sp. <i>N</i> is any integer: a number, a numeric register (r0.w.low), or a reference to a byte, word, or lword variable. The based indirect modes can take negative offsets as in B(p-8).		

these alternate addressing modes as a way to specify an explicit type in the variable reference. An &W(&p) tells you that p points at a word, information that you would not have if p were used by itself.

Note that *p* may not be preceded by an ampersand if it references a register because registers don't have addresses. Also, the ampersand preceding the *p* in the examples in Table 6.6 is optional if *p* is the name of an array or function. Since some compilers print a warning if you use the ampersand in front of an array or function name, you should leave the ampersand off unless your compiler requires it.

Finally, note that C-style pointer arithmetic is used when a double-indirect directive tells the assembler the type of the referenced object.

```
ptr p;
p += 1           /* Adds 1 to p because the referenced */
                  /*      object's type is unspecified. */
BP(&p) += 1     /* Adds 1 (size of a byte) to p. */
WP(&p) += 1     /* Adds 2 (size of a word) to p. */
LP(&p) += 1     /* Adds 4 (size of a lword) to p. */
```

The macros that implement the various type directives are in Listing 6.9. The prefixes themselves are declared, for the benefit of the compiler that's generating the C-code, in *c-code.h* and are shown in Listing 6.10.

Table 6.6. Combined Indirect and Effective-Address Modes

Syntax:	Evaluates to:
&p &W(&p) &WP(&p) &WP(fp+n)	address of the pointer
p &W(p) WP(&p) WP(fp+n)	contents of the pointer itself
W(p) *WP(&p) *WP(fp+n)	contents of the word whose address is in the pointer

Listing 6.9. *virtual.h*—Direct-Stack-Access Directives

```

45 #define W      * (word   *)
46 #define B      * (byte   *)
47 #define L      * (lword  *)
48 #define P      * (ptr    *)
49 #define WP     * (word   **)
50 #define BP     * (byte   **)
51 #define LP     * (lword  **)
52 #define PP     * (ptr    **)

```

Listing 6.10. *c-code.h*—Indirect-Mode Prefixes

```

12 #define BYTE_PREFIX      "B"      /* Indirect-mode prefixes. */
13 #define WORD_PREFIX       "W"
14 #define LWORD_PREFIX      "L"
15 #define PTR_PREFIX        "P"
16 #define BYTEPTR_PREFIX   "BP"
17 #define WORDPTR_PREFIX   "WP"
18 #define LWORDPTR_PREFIX  "LP"
19 #define PTRPTR_PREFIX    "PP"

```

6.2.7 Manipulating the Stack

Two C-code directives are provided to do explicit stack manipulation. These are:

Push and pop directives.

`push(something)`

and

`something = pop(type)`

The `push()` macro pushes the indicated object, and the `pop()` macro pops an object of the indicated type. For example, if `x` were declared as an `lword`, you could say: `x = pop(lword)`. The `lword` is the declared type of the target (of `x`). The target can also be a register, but the types on the two sides of the equal sign must agree:

Push junk if pushed object is too short.

```
rl.w.low = pop( word );
```

The stack is 32 bits wide, so part of the stack word is wasted when you push or pop small objects. For example:

r1.w.low = pop(word)

pushes the bottom word of the current top-of-stack item into the low half of $r1$. The top half of the stack item is discarded, and the top half of $r1$ is not modified. A push instruction modifies a 32-bit quantity, and the pushed object is right-adjusted in the 32-bit word. If $r1$ holds 0x12345678, the following instruction:

```
push( r1.b.b3 )
```

pushes the number 0x??????12. The question marks represent undefined values. A

```
push( r1.w.low )
```

directive pushes 0x????5678,

```
push( r1.w.high )
```

pushes 0x????1234, and

push(r1.w.l)

pushes the entire number 0x12345678.

A simple `push(r1)` is not permitted. The register name must be fully qualified. Use `push(r1.l)` to push the entire register.

The two stack directives are defined in *virtual.h* in Listing 6.11.

Listing 6.11. *virtual.h*—Pushing and Popping

```
53 #define push(n)          (--sp)->1 = (lword)(n)
54 #define pop(t)           (t) ( (sp++)->1 )
```

6.2.8 Subroutines

Subroutine definitions.

Subroutines must all be defined in the *text* segment. External-subroutine declarations should be output as follows:

external name();

No return value may be specified.

Subroutine definitions are created by surrounding the code that comprises the subroutine with PROC (name, class) and ENDP (name) directives. The name is the function name and the class is the storage class (either public or private). private functions cannot be accessed outside the current file. Invocations of PROC() and ENDP() may not be followed by semicolons.

Passing control:
call(name) **and** ret()

A subroutine is called using a `call (name)` directive and control is passed back to the calling function with a `ret()` directive. For example:

```
SEG ( text )
PROC( _sylvester, public )
    call( _tweety );           /* Call subroutine tweety.      */
    ret();                     /* Return to calling function. */
ENDP( sylvester )
```

The argument to `call()` can be either a subroutine name or a reference to a variable or register [like `call(r1.pp)`, which calls the subroutine whose address is in `r1.ppl`].

Note that the compiler has inserted underscores in front of the variable names so that these names won't conflict with internally generated labels. A `ret()` statement is supplied if one is not already present immediately before the `ENDP`.

The `virtual.h` definitions for these macros are in Listing 6.12. The `call()` macro simulates an assembly-language `call` instruction which pushes the return address (the address of the instruction that follows the call) onto the stack and then transfers control to the address that is the call's argument. Here, the return-address push is simulated by pushing the “stringized” version of the subroutine name onto the stack, and then the subroutine is called.⁷ An assembly-language return instruction pops the address at the top of stack into the instruction pointer. The `ret()` directive simulates this process by popping the name from the stack and returning.

Leading underscores avoid name conflicts.

`Virtual.h` implementation of `call()` and `ret()`.

Listing 6.12. `virtual.h`— Subroutine Definitions, Calls, and Returns

```

55 #define PROC(name,cls)    cls name() {
56 #define ENDP(name)        ret(); }           /* Name is ignored. */
57
58 #define call(name)        (--__sp)->pp = #name, (*(_void (*)())(name))()
59
60 #define ret()             __sp++; return

```

The subroutine call on line 58 of Listing 6.12 is complicated by the fact that function pointers must be handled as well as explicit function names. The argument must first be cast into a pointer to a subroutine:

```
(void (*)()) (name)
```

and then can be called indirectly through the pointer. This cast does nothing if *name* is already a subroutine name, but it correctly converts register and variable references into indirect subroutine calls.

6.2.9 Stack Frames: Subroutine Arguments and Automatic Variables

Of the various memory segments, the stack segment is of particular importance to subroutines. The run-time stack is used to make subroutine calls in the normal way: The return address is pushed as part of the call and the return pops the address at top of stack into the instruction pointer. Languages like C, which support recursion, use the stack for other purposes as well. In particular, the stack is used to pass arguments to subroutines, and certain local variables, called *automatic* variables, are stored on the stack at run time. In C, all local variables that aren't declared `static` are automatic. The following sequence of events occurs when a subroutine is called:

Subroutine linkage.

- (1) The arguments are pushed in reverse order.
- (2) The subroutine is called, pushing the return address as part of the call.
- (3) The called subroutine pushes a few housekeeping registers, including the frame pointer, discussed below.
- (4) The subroutine advances the stack pointer so that room is freed on the stack for local, automatic variables and anonymous temporaries.

Automatic variables.

Subroutine linkage: creating the stack frame.

7. The `#` directive is new to ANSI C. It turns the associated macro argument into a string by surrounding the matching text with implicit quotation marks.

Stack frame (activation record).

Building a stack frame.

This entire area of the stack, extending from the first argument pushed to the uppermost local variable or temporary, is called an *activation record* or *stack frame*. The compiler uses the stack and stack-manipulation macros described earlier to translate the code in Listing 6.13 into the output code in Listing 6.14. The complete stack frame is pictured in Figure 6.4.

Listing 6.13. call(of, the, wild): Compiler Input

```

1 int dog1, dog2, dog3;
2
3 static call( of, the, wild )
4 int of, the, wild;
5 {
6     int buck, thornton;
7     long john_silver;
8     ...
9 }
10 spitz()
11 {
12     call( dog1, dog2, dog3 );
13 }
```

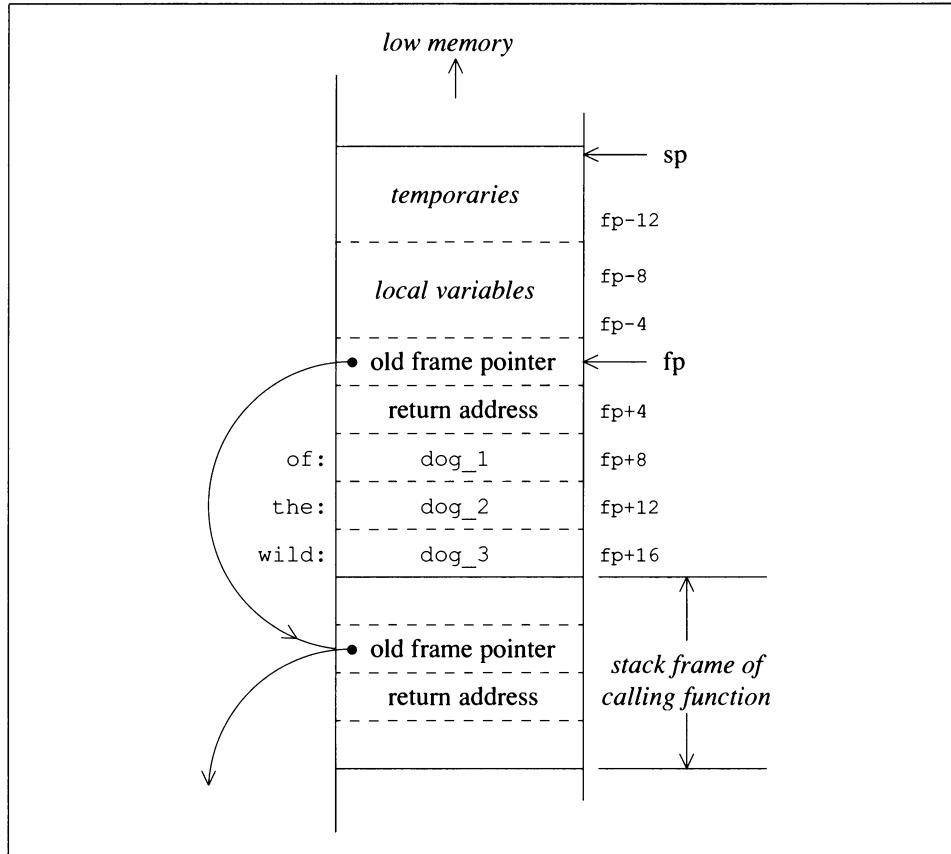
Listing 6.14. call(of, the, wild): Compiler Output

```

1 SEG( bss )
2 int      _dog1;
3 int      _dog2;
4 int      _dog3;
5
6 SEG (text )
7 PROC(_call, private)
8     push( fp );           /* Save old frame pointer. */
9     fp = sp;              /* Set up new frame pointer. */
10    sp -= 8;               /* Make room for local vars. & temporaries. */
11    call( chkstk );       /* Check for stack overflow. */
12    ...                   /* Code goes here. */
13    sp = fp;              /* Discard local variables and temporaries. */
14    fp = pop( ptr );      /* Restore previous subroutine's fp. */
15    ret();                /* */
16 ENDP(_call)
17 PROC(_spitz, public)
18     push( fp );          /* No sp-=N because there're no local vars. */
19     fp = sp;              /* */
20     call( chkstk );
21     push( _dog3 );        /* call( dog1, dog2, dog3 ); */
22     push( _dog2 );
23     push( _dog1 );
24     call( _call );
25     sp += 8;              /* Discard the arguments to call(). */
26     sp = fp;              /* Return from spitz(): discard local vars. */
27     fp = pop( ptr );      /* Restore calling routine's frame pointer. */
28     ret();                /* */
29 ENDP( _spitz )
```

Local variables in the stack frame.

The `sp-=8` on line ten of Listing 6.14 decrements the stack pointer to make room both for local variables and for a small scratch space to use for the anonymous temporaries. (Note that the stack pointer must be decremented in even multiples of 4 [the stack width]

Figure 6.4. Stack Frame for `call(of, the, wild)`

when you modify it explicitly, as compared to modifying it implicitly with a `push` or `pop`.) Unlike the arguments, the positions of local variables and temporaries within their area are arbitrary.⁸ Don't be confused by the fact that part of the stack frame is created by the calling function and the other part is created by the called function. The stack frame is a single entity. Everybody is responsible for undoing what they did, so the calling function cleans the arguments off the stack, and the called function gets rid of everything else. Note that this organization of the stack frame, though characteristic, is not mandatory. Many compilers put the temporaries in different places or not on the stack at all. Sometimes, space on the stack frame is reserved to pass arguments to run-time library functions, to push registers used for `register` variables, and so on.

Figure 6.5 shows the local-variable and argument portion of the stack frame in greater detail. The *frame-pointer* register (*fp*) provides a fixed reference into the stack frame. It is used to access all the arguments and automatic variables. The argument `wild`—which contains `dog_3`—can be accessed using the indirect addressing mode as follows:

The frame pointer: arguments and automatic variables.

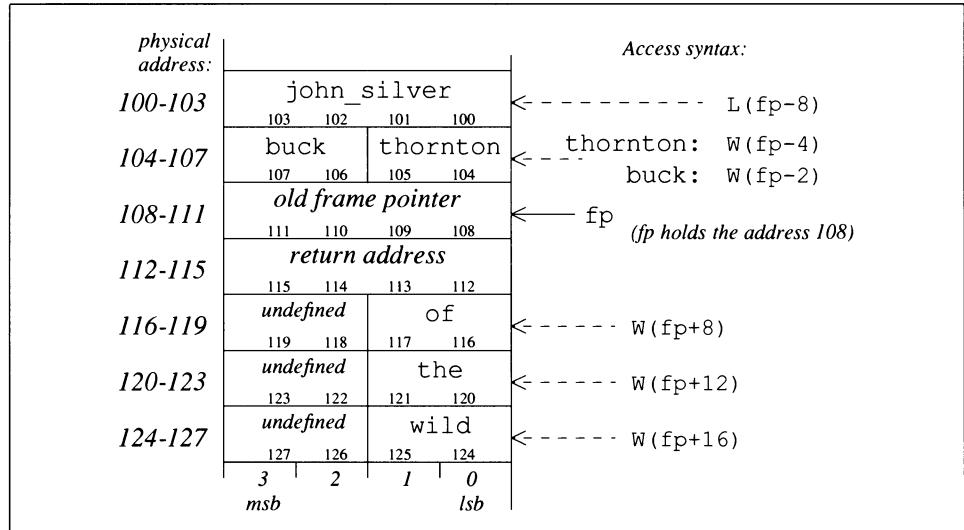
Accessing subroutine arguments.

8. The `chkstk()` subroutine, called on line 11, checks that the stack pointer hasn't crossed over into a different segment, as can happen if you allocate a large automatic array. It aborts the program if an error is discovered. This routine can also be used by a profiler subroutine to log the time at which the subroutine was entered. There would be a second subroutine call that logs the exit time at the bottom of the function.

$W(fp+16)$

The arguments are put onto the stack with push directives, and since an **int** is word sized, the high half of the 32-bit stack item is undefined. The actual, symbolic names of the argument do not appear in the output. All generated code that references `wild` uses $W(fp+16)$ instead of the symbolic name, `wild`.

Figure 6.5. The Stack Frame, Magnified



Accessing local automatic variables.

The local-variable region differs from the arguments in that the compiler treats a block of memory on the stack as a region of normal memory, and packs the local variables into that space as best it can. Again, the indirect modes are the best way to access these variables. For example, `buck` can be accessed using $W(fp-2)$. (`buck` is at physical addresses 106 and 107, and the frame pointer holds physical address 108, so the byte offset from one to the other is -2.) The $W(fp-2)$ directive causes a word-size object to be fetched from address 106. The low byte of `buck` can be accessed directly with $B(fp-2)$ and the high byte with $B(fp-1)$. Similarly, `john_silver` can be accessed with $L(fp-8)$. (Remember that the object's address is the physical address of the least-significant byte, so the offset is -8 here.) A pointer-to-word variable can be fetched with $*WP(fp-N)$. The object to which it points can be fetched with $*WP(fp-N)$.

Stack frame advantages:
size,

recursion.

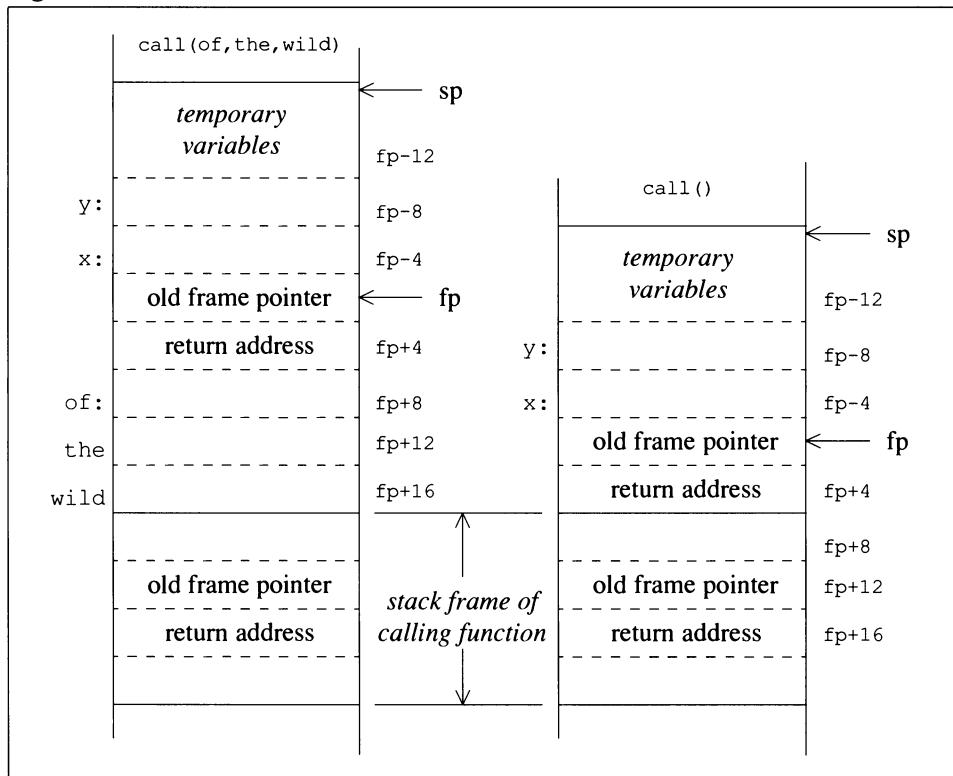
This use of the stack has two real advantages. First, the same relatively small area of memory (the stack) can be recycled from subroutine to subroutine, so less of the total memory area need be allocated for variables. Second, this organization makes recursive subroutines possible because each recursive instance of a subroutine has its own stack frame with its own set of local variables, and these variables are accessed relative to the current frame pointer. A recursive subroutine does not know that it's calling itself; it does the same thing for a recursive call that it would do for a nonrecursive call: push the arguments and transfer control to the top of the required subroutine. The called routine doesn't know that it has been called recursively, it just sets up the stack frame in the normal way. It doesn't matter if more than one stack frame for a given subroutine exists at once—only the top one is active.

Also note that the address of the leftmost argument is always $(fp+8)$, regardless of the number of arguments or their type. This is one of the things that makes it possible to have a variable number of arguments in a C subroutine—you can always find the leftmost one.

The stack frame's organization has disadvantages, too. The problem is that the code that pushes the arguments is generated by the compiler when it processes the subroutine call, but the offsets are figured when the compiler processes the subroutine declaration. Since the call and declaration can be in different files, there's no way that the compiler can check for consistency unless you use a function prototype in a common header file. If you don't use the prototype, a particularly nasty bug, called a *phase error*, can appear at run time. Figure 6.6 shows the stack frames created, both by the earlier `call(of, the, wild)` and an incorrect `call()`, with no arguments. When the `call()` subroutine modifies `wild`, it just modifies the memory location at `fp+4`, and on the incorrect stack, ends up modifying the return address of the calling function. This means that `call()` could work correctly, as could the calling function, but the program would blow up when the calling function returned.

Stack frame disadvantages: phase errors.

Figure 6.6. A Phase Error



The sequence of instructions that set up a stack frame are so common that C-code provides two instructions for this purpose. The `link(N)` instruction does the following:

```
push( fp );
fp = sp;
sp -= N * stack_width; /* Decrement by one stack element. */
```

and an `unlink()` directive does the following:

```
sp = fp;
fp = pop( ptr );
```

link and *unlink* instructions.

The earlier `call(of, the, wild)` is modified to use `link()` and `unlink()` in Listing 6.15. Listing 6.16 shows the `link` and `unlink` implementations in `virtual.h`.

Listing 6.15. call(of, the, wild): Compiler Output with link() and unlink()

```

1 SEG( bss )
2 int      _of;
3 int      _the;
4 int      _wild;
5
6 SEG (text )
7 PROC(_call)
8     link( 2 );      /* Create stack frame, 2 stack elements for locals. */
9     ...
10    unlink();
11    ret();
12 ENDP(_call)
13
14 PROC(_spitz)
15     link( 0 );      /* No local variables so argument is zero.          */
16     push( _wild );  /* call( of, the, wild );                                */
17     push( _the );
18     push( _of );
19     call( _call );
20     unlink();
21     ret();
22 ENDP(_spitz)

```

Listing 6.16. virtual.h—Subroutine Linkage Directives

```

61 #define link(n)   ((--__sp)->pp = (char *)__fp) , (__fp = __sp) , (__sp -= (n))
62
63 #define unlink()   (__sp = (reg *)__fp) , (__fp = (reg *)((__sp++)->pp))

```

Dynamic, static links.

The stack frames can also be viewed as a linked list of data structures, each representing an active subroutine. It's possible (in an object-oriented programming environment, for example) for these structures to be allocated from discontinuous memory rather than from a stack. The frame pointer is sometimes called the *dynamic link* because it links together activation records.⁹

6.2.10 Subroutine Return Values

A function's return values should be placed into a register, according to Table 6.7.

9. Some languages require a second pointer, called the *static link*, discussed in Appendix B.

Table 6.7. Return Values: Register Usage

Type	Returned in:
char	rF.w.low (A char is always promoted to int .)
int	rF.w.low
long	rF.l
pointer	rF.pp

6.2.11 Operators

A very limited set of operators are supported in C-code. The following arithmetic operators (only) are available:

```
= + = - = * = / = %= ^= &= |= <<= >>= == =~ lrs(x,n)
```

Arithmetic operators.

The ones that look like C operators work just like the equivalent C operators. $>>=$ is an arithmetic right shift (sign extension is assumed). Three special operators are supported. The $=-$ operator performs a two's complement operation, as in $x =- c$. The $=\sim$ operator does a one's complement, as in $x =\sim c$. Both of these instructions are perfectly legitimate C. I'm assuming that the lexical analyzer breaks the foregoing into:

```
x = -c ;
x = \sim c ;
```

Lint might print an "obsolete syntax" error message, though. You need binary operators here, because C-code is attempting to mimic real assembly language, and though most machines support a unary "negate" operator such as the 8086 NEG AX, which does a two's-complement negation of the AX register, you can't implement this operator without an equal sign in C. It seemed better to use a syntax consistent with the other operators than to bury the operation in an macro.

One additional operation is supported by means of a macro. A `lrs(x,n)` directive does a logical right shift rather than an arithmetic shift ($>>=$). It shifts `x`, `n` bits to the right, with zero fill in the high bits rather than sign extension. This directive is defined in `virtual.h` in Listing 6.17.

Listing 6.17. `virtual.h`—The Logical-Right-Shift Directive

```
64 #define lrs(x,n) ((x) = ((unsigned long)(x) >> (n)))
```

All C-code expressions must contain exactly two operands and one operator. They must be semicolon terminated. Only one expression is permitted on a line, and the entire expression must be on a single line.

6.2.12 Type Conversions

No automatic type conversions are supported in C-code. Both operands in an expression must be of the same type, but if one of the operands is an explicit (immediate-mode) number, it is automatically converted to the type of the other operand.¹⁰

10. Note that, because of this restriction, most real assembly languages attach the type to the operator rather than the operand. You would move a 32-bit long word in 68000 assembler with a `MOV.L d1,d0` instruction. In C-code you'd use `r1.l = r0.l` to do the same thing.

Sign extension.

Cast operators may not be used to do type conversion, but sign extension can be performed using one of the following directives:

<code>ext_low(reg)</code>	Duplicate high bit of <code>reg.b.b0</code> in all bits of <code>reg.b.b1</code>
<code>ext_high(reg)</code>	Duplicate high bit of <code>reg.b.b2</code> in all bits of <code>reg.b.b3</code>
<code>ext_word(reg)</code>	Duplicate high bit of <code>reg.w.low</code> in all bits of <code>reg.w.high</code>

An `ext_low(reg)` directive fills `reg.b.b1` with ones if the sign bit in `reg.b.b0` is set, otherwise it's filled with zeros. These directives only work in a register, so given input like this:

```
int archy;
long mehitabel;
mehitabel = archy + 1;
```

the compiler should output something like the following:

```
public word _archy;
public lword _mehitabel;

r0.w.low = _archy; /* Get archy. */
r0.w.low += 1; /* Add 1 to it. */
ext_word( r0 ); /* Convert to long. */
_mehitabel = r0.l; /* Do the assignment. */
```

The `ext_word()` directive effectively converts the word into an lword. If `archy` were `unsigned`, a `r0.w.high=0` would be used instead of `ext_word(r0)`. The definitions of the sign-extension directives are in Listing 6.18.

Listing 6.18. `virtual.h`—Sign-Extension Directives

```
65 #define ext_low(reg) (reg.w.low = (word )reg.b.b0 )
66 #define ext_high(reg) (reg.w.high = (word )reg.b.b2 )
67 #define ext_word(reg) (reg.l = (lword)reg.w.low )
```

6.2.13 Labels and Control Flow

Labels and the `goto`.

Labels in C-code are like normal C labels. They are defined as a name followed by a colon:

```
label:
```

The only control-flow statement is the `goto` branch, which is used in the normal way:

```
goto label;
```

The target of the `goto` branch must be in the same subroutine as the `goto` itself.

Test directives.

Conditional flow of control is performed using one of the test directives summarized in Table 6.8. These all compare two operands, and the instruction on the line following the test is executed only if the test evaluates true, otherwise the instruction on the next line is ignored. The following code executes the `goto` branch if `all` is equal to `things` (all things being equal).

```
EQ( all, things )
goto jail;
go:
/* collect $200 */
jail:
```

All the addressing modes described earlier can be used in a test. The normal

comparisons assume signed numbers, but the `U_LT()`, `U_LE()`, `U_GT()`, and `U_GE()` directives compare the two numbers as unsigned quantities. The instruction following the test may not be another test. The test directives are implemented with the macros in Listing 6.19.

Table 6.8. C-code Test Directives

Directive:	Execute following line if:
<code>EQ(a, b)</code>	$a = b$
<code>NE(a, b)</code>	$a \neq b$
<code>LT(a, b)</code>	$a < b$
<code>LE(a, b)</code>	$a \leq b$
<code>GT(a, b)</code>	$a > b$
<code>GE(a, b)</code>	$a \geq b$
<code>U_LT(a, b)</code>	$a < b$ (unsigned comparison)
<code>U_LE(a, b)</code>	$a \leq b$ (unsigned comparison)
<code>U_GT(a, b)</code>	$a > b$ (unsigned comparison)
<code>U_GE(a, b)</code>	$a \geq b$ (unsigned comparison)
<code>BIT(b, s)</code>	bit b of s is set to 1 (bit 0 is the low bit).

Listing 6.19. `virtual.h`—Comparison Directives

```

68 #define EQ(a,b)           if( (long)(a) == (long)(b) )
69 #define NE(a,b)           if( (long)(a) != (long)(b) )
70 #define LT(a,b)           if( (long)(a) <  (long)(b) )
71 #define LE(a,b)           if( (long)(a) <= (long)(b) )
72 #define GT(a,b)           if( (long)(a) >  (long)(b) )
73 #define GE(a,b)           if( (long)(a) >= (long)(b) )
74
75 #define U_LT(a,b)          if( (unsigned long)(a) <  (unsigned long)(b) )
76 #define U_GT(a,b)          if( (unsigned long)(a) >  (unsigned long)(b) )
77 #define U_LE(a,b)          if( (unsigned long)(a) <= (unsigned long)(b) )
78 #define U_GE(a,b)          if( (unsigned long)(a) >= (unsigned long)(b) )
79
80 #define BIT(b,s)           if( (s) & (1 << (b)) )

```

6.2.14 Macros and Constant Expressions

Various C-preprocessor macros, summarized in Table 6.9, may be used in a C-code file. C-style, parameterized macros are supported, but the ANSI `defined()` pseudo op, the concatenation operator (`##`), and the “stringizing” operator (`#name`) are not supported. Arithmetic expressions that involve nothing but constants and the following operators can appear in an `#if` directive and anywhere that a constant would appear in an operand:

+ - * / % & && | || == != < > <= >= ~ !

The `-` is both unary and binary minus. The `*` is multiplication. Multiple-operator expressions such as the following are legal:

```
#define L1 (-24)
WP(fp-4) = WP(fp-L1-6); /* The L1-6 is a constant expression. */
```

Table 6.9. C-code Preprocessor Directives

#line <i>line-number</i> "file"
#define NAME <i>text</i>
#define NAME (<i>args</i>) <i>text</i>
#undef NAME
#ifdef NAME
#if <i>constant expression</i>
#endif
#else
#include <file>
#include "file"

Note that the constant expression must completely precede or follow a variable or register reference in an operand. For example, the following is not permitted because the register reference (fp) is imbedded in the middle of the expression:

```
WP(fp-4) = WP(6 + fp + 10);
```

6.2.15 File Organization

All variable and macro declarations must precede their use. Forward references are permitted with **goto** branches, however.

C-code files should, ideally, take the following form:

```
#include <tools/virtual.h>
SEG( data )
    initialized data declarations
SEG( bss )
    uninitialized data declarations
SEG( text )
    subroutines
```

Can't issue SEG() between PROC() and ENDP().

You can switch back and forth between the *data* and *bss* segments as much as you like, but you can't change segments in the middle of a subroutine definition bounded by a PROC() and ENDP() directive. **extern** declarations should be placed in the *bss* segment. Macro definitions can go anywhere in the file.

6.2.16 Miscellany

Translate main() to main().

One other handy macro is in *virtual.h*. Since the point of C-code is to create a language that can be assembled with your C compiler rather than a real assembler, it's desirable to be able for a *main()* function in the source to really be *main()*. Unfortunately, the compiler puts a leading underscore in front of *main* when it outputs the PROC directive. The main() macro on line 81 of Listing 6.20 takes care of this problem.

Print virtual-machine state, pm().

The last part of *virtual.h* (on lines 84 to 118 of Listing 6.20) is an actual function, created only if ALLOC is defined. *pm()* prints the top few stack elements and the contents of the virtual-machine registers to standard output as follows:

Listing 6.20. *virtual.h*—Run-time Trace Support

```

81 #define _main main
82
83 #ifdef ALLOC
84 pm()
85 {
86     reg *p;
87     int i;
88
89     /* Print the virtual machine (registers and top 16 stack elements). */
90
91     printf("r0= %08lx r1= %08lx r2= %08lx r3= %08lx\n",
92           r0.l, r1.l, r2.l, r3.l);
93     printf("r4= %08lx r5= %08lx r6= %08lx r7= %08lx\n",
94           r4.l, r5.l, r6.l, r7.l);
95     printf("r8= %08lx r9= %08lx rA= %08lx rB= %08lx\n",
96           r8.l, r9.l, rA.l, rB.l);
97     printf("rC= %08lx rD= %08lx rE= %08lx rF= %08lx\n",
98           rC.l, rD.l, rE.l, rF.l);
99
100    if( __sp >= &stack[SDEPTH] )
101        printf("Stack is empty\n");
102    else
103        printf("\nitem byte real addr b3 b2 b1 b0 hi lo 1\n");
104
105    for( p = __sp, i=16; p < &stack[SDEPTH] && --i>=0; ++p )
106    {
107        printf("%04d %04d %9p [%02x|%02x|%02x|%02x] = [%04x|%04x] = [%08lx]",
108              p-__sp, (p-__sp)*4, (void far *)p,
109              p->b.b3 & 0xff, p->b.b2 & 0xff, p->b.b1 & 0xff, p->b.b0 & 0xff,
110              p->w.high & 0xffff, p->w.low & 0xffff,
111              p->l);
112    }
113
114    if( p == __sp ) printf("<-SP");
115    if( p == __fp ) printf("<-FP");
116    printf("\n");
117 }
118 }
119#endif

```

```

r0= 00000000 r1= 00000000 r2= 00000000 r3= 00000000
r4= 00000000 r5= 00000000 r6= 00000000 r7= 00000000
r8= 00000000 r9= 00000000 rA= 00000000 rB= 00000000
rC= 00000000 rD= 00000000 rE= 00000000 rF= 00000000

item byte real addr b3 b2 b1 b0 hi lo 1
0000 0000 2C8F:1878 [00|00|00|00] = [0000|0000] = [00000000]<-SP
0001 0004 2C8F:187C [00|00|00|00] = [0000|0000] = [00000000]
0002 0008 2C8F:1880 [00|00|18|8c] = [0000|188c] = [0000188c]<-FP
0003 0012 2C8F:1884 [00|00|04|f2] = [0000|04f2] = [000004f2]
0004 0016 2C8F:1888 [ab|cd|ef|12] = [abcd|ef12] = [abcdef12]
0005 0020 2C8F:188C [00|00|18|90] = [0000|1890] = [00001890]

```

The stack is printed three times, broken up as bytes, words, and 1words. The leftmost, *item*, column is the offset in stack elements from the top of stack to the current element; the *byte* column is the byte offset to the least-significant byte of the stack item; and the *real addr* is the physical address of the stack element (it's in 8086 segment:offset form).

The position of the `sp` and `fp` register is indicated at the far right. `pm()` is used in the compiler to print a run-time trace as the output code executes. The compiler described below can be put in a mode where most instructions are output like this:

```
r0.1 = r1.1;           printf("r0.1 = r1.1;\n"); pm();
```

so you can what the virtual machine is doing as the code executes.

C-code labels are different from assembly language.

Jump tables can't be implemented in C-code.

6.2.17 Caveats

C-code is similar enough to assembly language that it's easy to forget it's really C, and make mistakes accordingly. The biggest problem is labels. Most assembly languages treat all labels the same, regardless of what they are used for: a label always evaluates to an address, and any label can be used in any instruction. C-code is different, however: Array names evaluate to addresses, most other labels evaluate to the contents of a variable, but labels that are followed by colons can only be used as targets of `goto` branches. These restrictions make C-code a little more difficult to write than real assembler. They also make some common assembly-language data structures impossible to implement. For example, most assemblers let you code a `switch` statement as follows:

```
SEG( data )
switch_table: dw L1           /* Array of four labels. */
               dw L2
               dw L3
               dw L4
SEG( text )

t0 = /* evaluated argument to switch */
goto switch_table[ t0 ]
L1:   code...
L2:   code...
L3:   code...
L4:   code...
```

This data structure, called a *jump table*, is not legal C-code—you can't have an array of labels and you can't use a variable as an argument to a `goto`.

Subroutine prefix, body, and suffix.

The symbol table.

6.3 The Symbol Table

Like Gaul, most computer languages can be divided into three parts: declarations, expressions, and statements. Looking back at Table 6.2 on page 451, the output code is divided into three logical sections. Each section is generated from a separate part of the input file, and the code generation for each of these parts is controlled by distinct parts of the grammar. Every subroutine has a *prefix* portion that's generated when the declarations and argument list are processed. The *body* of the subroutine, which contains statements and expressions, comes next. Finally, code to clean up and return from a subroutine is generated at the end of the definition. This last section is called the *suffix*. I'll start by looking at the first of these sections, the subroutine-prefix and declaration processing.

6.3.1 Symbol-Table Requirements

We need to examine the data structures that are used to process a declaration before looking at the actual code-generation actions. A compiler's declaration system centers around a set of data structures collectively called the *symbol table*. Strictly speaking, the

symbol table is a database that contains information about subroutines, variables, and so forth. The database is indexed by a *key* field—here a subroutine or variable’s name—and each record (each entry in the database) contains information about that item such as the variable’s type or subroutine’s return value. A record is added to the database by the code that processes declarations, and it is deleted from the database when the scoping rules of the language determine that the object can no longer be referenced. C local variables, for example, are deleted when the compiler finishes the block in which they are declared.

Key, record.

Symbol tables are used for other purposes as well. Type definitions and constant declarations may be found in them, for example. The symbol table can also be used to communicate with the lexical analyzer. In the current compiler, a **typedef** creates a symbol-table entry for the new type, as if the type name were a variable name. A bit is set in the record to indicate that this is a **typedef**, however. The lexical analyzer then uses the symbol table to distinguish identifiers from type names. This approach has its drawbacks (discussed below), but can be quite useful.

Even though a symbol table is a database, it has special needs that must be met in specific ways. A symbol-table manager must have the following characteristics:

Desirable symbol-table characteristics.

- Speed. Because the symbol table must be accessed every time an identifier or type is referenced, look-up time must be as fast as possible. Consequently, disk-based data-management systems are not appropriate here. The entire table should be in memory. On the down side, one of the main limitations on input-file size is often the maximum memory available for the symbol table.
- Ease of maintenance. The symbol table is probably the most complex data structure in the compiler. Its support functions must be organized so that someone other than the compiler writer can maintain them.
- Flexibility. A language like C does not limit the complexity of a variable declaration, so the symbol table must be able to represent variables of arbitrary type. This representation should be optimized for code-generation purposes. Similarly, the symbol table should be able to grow as symbols are added to it.
- Duplicate entries must be supported. Most programming languages allow a variable at an inner nesting level to have the same name as a variable at an outer nesting level. These are different variables, in spite of having the same name, and the scoping rules of the language determine which of these variables are active at a given moment. The active variable is said to *shadow* the inactive one. A distinct symbol table entry is required for each variable, and the database manager must be able to handle this situation.
- You must be able to quickly delete arbitrary elements and groups of elements from the table. For example, you should be able to delete all the local variables associated with a particular block level in an efficient manner, without having to look up each element separately.

Shadowing.

The symbol table used here is organized in two layers. I’ll call the innermost of these the *database* layer. This layer takes care of physical table maintenance: inserting new entries in the table, finding them, deleting them, and so forth. I’ll call the outer layer the *maintenance* layer—it manages the table at a higher level, creating systems of data structures to represent specific symbols and inserting these structures into the table using the low-level insert function. Other subroutines at the maintenance level delete entire classes of symbols (variables declared at the same block level, for example), traverse all variables of a single class, and so forth.

Layers
Database layer.

Maintenance layer.

Stack-based symbol tables.

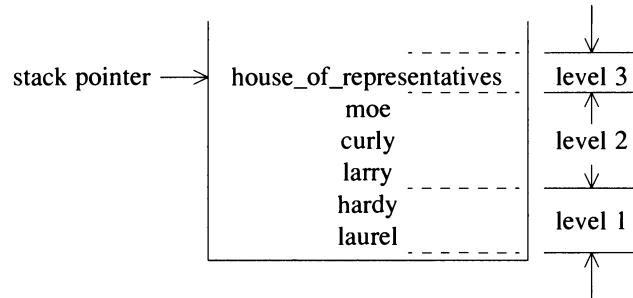
6.3.2 Symbol-Table Data-Base Data Structures

Several data structures can be used for the database layer of the symbol table, each appropriate in specific situations. The simplest possible structure is a linear array organized as a stack. New symbols are added to the end of the array with a push operation and the array is searched from top to bottom of stack. (The most recently added item is examined first.) This method, though crude, is quite workable provided that the table is small enough. The scope rules are handled by the back-to-front searching. Since variables declared at an inner nesting level are added to the table after those declared at a higher level, they are always found first. The database manager is trivial to implement.

One real advantage to a stack-based symbol table is that it's very easy to delete a block of declarations. Variable declarations are done in waves according to the current scoping level. For example, given input like this:

```
int laurel, hardy;
{
    int larry, curly, moe;
    {
        int house_of_representatives[ 435 ];
    }
}
```

laurel and hardy are inserted first as a block, then larry, curly, and moe are inserted, and then the house_of_representatives is inserted. The symbol-table stack looks like this when the innermost block is processed:



All variables associated with a block can be deleted at one time by adding a constant to the stack pointer.

Stack disadvantages.

The stack approach does have disadvantages other than the obvious one of the inefficient linear search required to find entries close to the beginning of the table. (This is not an insignificant problem—the compiler spends more time searching for symbol-table references than any other symbol-table operation. The time required for a linear search will be prohibitive if the table is large enough.) The maximum size of the stack-based table must be known at compile time, so the symbol table can't scale its size dynamically to fit the needs of the current input file. Consequently, the number of variables that can be handled by the system is limited. You have to allocate a worst-case table size to make sure that there is enough room.

Tree-based symbol tables.

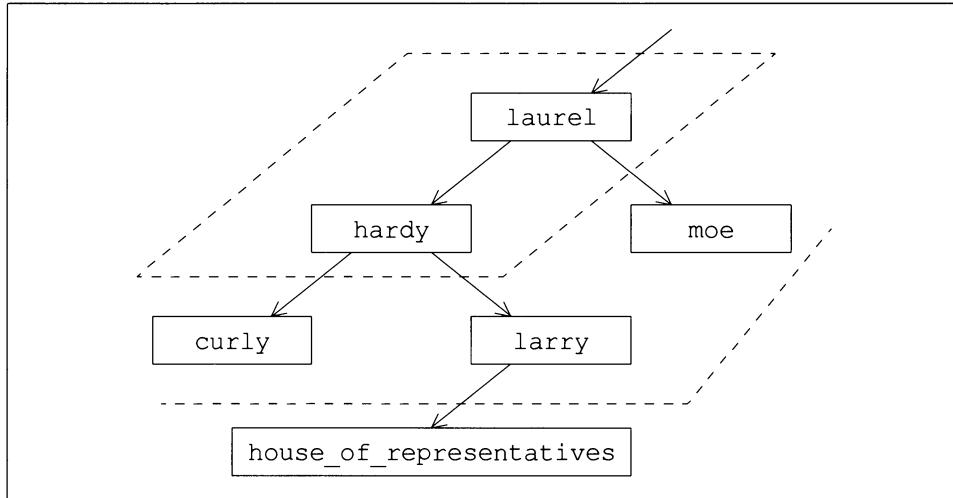
Tree deletions.

The search-time and limited-size problems can be solved, to some extent, by using a binary tree as the basic data structure. Average search times in a balanced binary tree are logarithmic, and the tree size can grow dynamically as necessary.

Deletion of an arbitrary node from a tree is difficult, but fortunately this is not an issue with most symbol-tables applications, because nodes for a given level are inserted into the tree as a block, and newer levels are deleted before the older ones. The most-recently inserted nodes tend to form leaves in the tree. If a variable at the most recent level is an interior node, then all its children will have been declared either at the same

nesting level or in an inner block. The most recently added block of variables is always at the end of a branch, and all variables in that block can be removed by breaking the links to them without having to rearrange the tree. For example, a tree for the earlier code fragment is pictured in Figure 6.7. The dotted lines show the scoping levels. The `house_of_representatives` is deleted first by breaking a single link; `larry`, `curly`, and `moe` are deleted next, again by breaking single links; `laurel` and `hardy` are deleted by breaking the single link that points at `laurel`.

Figure 6.7. Symbol-Table Trees: Deletion



Binary trees do have disadvantages. First, it is a common practice for programmers to declare variables in alphabetical order. Since variables are added to the tree in the same order as the declarations, a simple binary tree degrades to a linked list in this situation, and the search times are linear rather than logarithmic. This problem can be solved at the cost of greater insert and delete times by using a height-balanced tree system such as an AVL tree,¹¹ but the shuffling around of nodes that's implicit in the rebalancing can destroy the ordering that made deletions easy to do. The other disadvantage of a tree-based symbol table is that *collisions*, situations where the same name is used for variables at different scope levels, are difficult to resolve. This problem, too, can be solved at the expense of lookup time. For example, you could have two key fields in the data structure used to represent a tree node, one for the name and one for the nesting level:

```

typedef struct tree_node
{
    char           name[ 32 ]; /* Variable name.          */
    int            level;      /* Nesting level.          */
    struct tree_node *right;    /* Right-child pointer.    */
    struct tree_node *left;     /* Left-child pointer.     */
    INFO_TYPE      info;       /* Other Information.     */
}
tree_node;
    
```

and then use both fields when comparing two nodes:

Tree disadvantages.

Collisions.

11. See: [Kruse] pp. 357–371 and [Tenenbaum] pp. 461–472 for a general discussion of AVL trees. [Holub1] contains a C implementation of an AVL-tree database manager.

```

compare( node1, node2 )
struct tree_node *node1, *node2;
{
    if( node1->level != node2->level )
        return( node1->level - node2->level );
    else
        return( strcmp(node1->name, node2->name) );
}

```

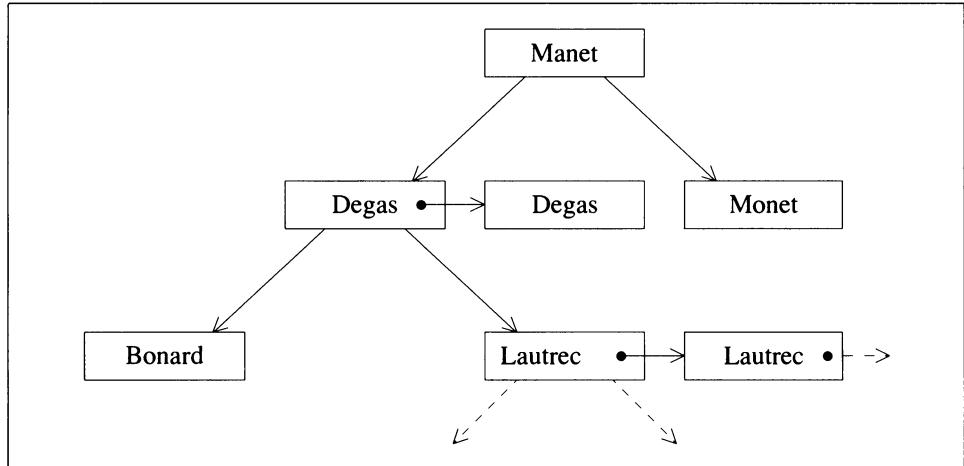
You could also solve the collision problem by adding an additional field to the tree node—a pointer to the head of a linked list of conflicting nodes. Newly added entries would be put at the head of the list, and so would be found first when the list was searched. The system is pictured in Figure 6.8; it uses the following data structure:

```

typedef struct tree_node
{
    char name[32];
    struct tree_node *right; /* Right-child pointer. */
    struct tree_node *left; /* Left-child pointer. */
    struct tree_node *conflicts; /* Conflicting-name-list head. */
    INFO_TYPE info; /* Other information. */
}
tree_node;

```

Figure 6.8. Using a Linked List to Resolve Collisions in a Tree-Based Table



There is one final problem with a tree-based symbol table. The use of global variables is discouraged by most proponents of structured-programming because it's difficult to determine how global variables change value. As a consequence, a well-structured program accesses local variables more often than global ones. The local variables are added to the symbol table last, however, and these nodes tend to be farther away from the root in a tree-based table, so it takes longer to find them.

Hashed symbol tables.

It turns out that the best data structure for most symbol-table applications is a *hash table*. An ideal hash table is an array that is indexed directly by the key field of the object in the table. For example, if the key field is a string made up of letters, digits, and underscores, the string could be treated as a base 63 number (26 lower-case letters + 26 upper-case letters + 10 digits + an underscore = 63 possible characters). Unfortunately, an array that could be indexed directly by a 16-character name would require $63^{16}-1$ or roughly 60,000,000,000,000,000,000,000 elements, most of which would be wasted because an average symbol table has only a few hundred objects in it.

The solution to this problem is to compress the array. A *hash table* is an array, indexed by key, that is compressed so that several elements of the uncompressed array can be found at a single location in the compressed array. To do this, you convert the key field that's used as the index in the uncompressed array into a pseudo-random number which is used as an index into the compressed array. This randomization process is called *hashing* and the number so generated is the key's *hash value*. The same key should always *hash* to the same pseudo-random number, but very similar keys should have very different hash values.

Hash tables.

Collisions (situations where two keys hash to the same value) are resolved by making each array element the head of a linked list of table elements.¹² This method is appropriate in a symbol-table application because, if you always put the new node at the head of a chain rather than the end, local variables automatically preempt global variables with the same name—they'll be found first when the list is searched. They are also found more quickly than they would be if stored in a binary tree.

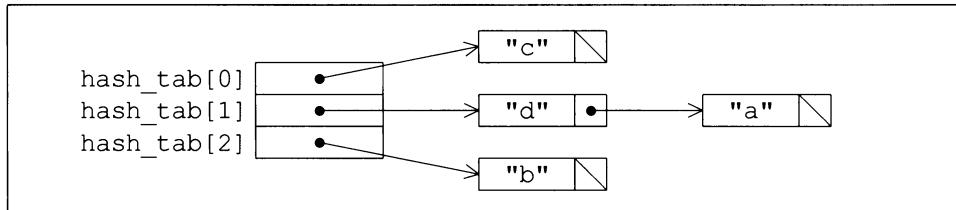
Hashing and hash values.

A simple hash table of size three is pictured in Figure 6.9. The table has four members, with the keys "a", "b", "c", and "d". The hash values are computed by treating the name strings as numbers. An ASCII 'a' has the decimal value 97, 'b' is 98, and so on. These numbers are truncated down to an array index using a modulus division by the table size (3 in this case):

Collision resolution.

name	numeric value	value MOD table size
"a"	97	1
"b"	98	2
"c"	99	0
"d"	100	1

Figure 6.9. A Simple Hash Table



A simple hash function could just add together the characters in the name as if they were numbers, and then truncate the resulting sum to a valid array index with a modulus operation. A better function is *hashpjw*, shown in Listing 6.21.¹³ *Hashpjw* uses an exclusive-or and shift strategy to randomize the hash value.

Two hash functions: addition and hashpjw.

There has probably been more waste paper devoted to the subject of hash algorithms than any other topic in Computer Science. In general, the complex algorithms so described give very good theoretical results, but are so slow as to be impractical in a real application. An ideal hash function generates a minimum number of collisions, and the

An informal complexity analysis of addition and hashpjw.

12. The collision-resolution method used here is called “open” hashing. There are other ways of resolving collisions that won't be discussed because they are not much use in most practical applications. See [Tenenbaum] pp. 521–574 and [Kruse] pp. 112–135.

13. Hashpjw was developed by P.J. Weinberger and is described in [Aho], p. 436. The version in Listing 6.21 is optimized somewhat from the one in [Aho] and is tailored to a 16-bit `int`. There's a more general purpose version of the same function in Appendix A.

Listing 6.21. *hashpjw.c*—A 16-bit *hashpjw* Implementation

```

1 unsigned hash_pjw( name )
2 unsigned char *name;
3 {
4     unsigned hash_val = 0;
5     unsigned i;
6
7     for(; *name ; ++name )
8     {
9         hash_val = (hash_val << 2) + *name ;
10        if( i = hash_val & 0x3fff )
11            hash_val = (hash_val ^ (i >> 12)) & ~0x3fff ;
12    }
13
14    return hash_val;
15 }
```

average search time (which is proportional to the mean chain length) should be roughly the same as an equivalent data structure such as a binary tree—it should be logarithmic. The maximum path length in a perfectly balanced tree with N elements is $\log_2 N$. A tree with 255 elements has a maximum chain length of 8 and an average chain length of roughly 7 [($128 \times 8 + 64 \times 7 + 32 \times 6 + \dots + 2 \times 1$)/255]. You're okay as long as you can get a comparable average chain length in a hash table.

It's difficult to predict chain lengths from a given algorithm because the words found in the input control these lengths. An empirical analysis is quite workable, however, given a large enough data set. Table 6.10 shows the behavior of the *hash_add()* and *hash_pjw()* functions from Appendix A. The former just adds together the characters in the name; the latter is in Listing 6.21. 917 unique variable names gleaned from real C programs were used for the test input, and the table is 127 elements long.

The elapsed time is the amount of time used to run the test program. Since the only change was the choice of hash function, this number gives us a good indication of efficiency. You'll note that *hash_pjw()* is about 9% slower than *hash_add()*. The mean chain length for both algorithms is identical, and both algorithms use shorter average chain lengths than a binary tree of an equivalent size. Finally, the distribution of chain lengths is actually a little better with *hash_add()* than with *hash_pjw()*—there are more shorter chains generated by *hash_add()*.

Addition as a hash function: advantages, disadvantages.

Weighing the foregoing, it seems as if simple addition is the preferable of the two algorithms because it yields somewhat shorter chains, executes marginally faster, and is smaller. Addition has two significant disadvantages, the main one being that characters are small numbers, so the hash values tend to pile up at one end of a large table. One of the reasons that addition performed well in the current example is the relatively small table size. *hash_pjw()* solves the bunching problem with the left shift on line nine of Listing 6.21, making it more appropriate for larger tables. Also, addition can't distinguish between identifiers that are permutations of the same name and *hashpjw* can, so the latter is a better choice if this situation comes up. Aho, using names gleaned from Pascal programs rather than C programs and a larger table size, got better results from *hashpjw* than I did.

Hash tables are almost ideal data structures for use in a symbol table. They are as efficient (if not more so) than binary trees, and are more appropriate for those applications in which several elements of the table might have the same key, as is the case when local and global variables share the same name.

Table 6.10. Performance of Two Hash Functions

Addition		Hashpjw	
Elapsed time	9.06 seconds	Elapsed time	9.61 seconds
Mean chain length:	7.22047	Mean chain length:	7.22047
Standard deviation:	2.37062	Standard deviation:	2.54901
Maximum chain length:	14	Maximum chain length:	13
Minimum chain length:	2	Minimum chain length:	1
0 chains of length	1	1 chain of length	1
1 chains of length	2	3 chains of length	2
8 chains of length	3	9 chains of length	3
7 chains of length	4	8 chains of length	4
18 chains of length	5	12 chains of length	5
13 chains of length	6	15 chains of length	6
22 chains of length	7	21 chains of length	7
22 chains of length	8	15 chains of length	8
16 chains of length	9	15 chains of length	9
8 chains of length	10	15 chains of length	10
7 chains of length	11	10 chains of length	11
2 chains of length	12	0 chains of length	12
2 chains of length	13	3 chains of length	13
1 chain of length	14	0 chains of length	14

6.3.3 Implementing the Symbol Table

I will use the hash functions described in Appendix A for the database layer of our symbol table. To summarize the appendix, the hash-table system described in Appendix A uses two data structures: the hash table itself is an array of pointers to “buckets”, each bucket being a single data-base record. The buckets are organized as a linked list. When two keys hash to the same value, the conflicting node is inserted at the beginning of the list. This list of buckets is doubly linked—each bucket contains a pointer to both its predecessor and successor in the list. Arbitrary nodes can then be deleted from the middle of a chain without having to traverse the entire chain. You should review these functions now.

The basic hash functions are fine for simple symbol-table applications. They were used to good effect in `occs` and `LLama`, for example. Most programming languages require a little more complexity, however. First of all, many internal operations require the compiler to treat all variables declared at a common nesting level as a single block. For example, local variables in C can be declared at the beginning of any curly-brace delimited compound statement. A function body is not a special case, it’s treated identically to a compound statement attached to a `while` statement, for example. The scope of any local variable is defined by the limits of the compound statement. For all practical purposes, a variable ceases to exist when the close curly-brace that ends that block in which it is declared is processed. The variable should be removed from the symbol table at that time. (Don’t confuse compile and run time here. A `static` local variable continues to exist at run time, but it can be deleted from the symbol table because it can’t be accessed at compile time once it is out of scope.)

Block deletions.

The compiler also has to be able to traverse the list of local variables in the order that they were declared. For example, when the compiler is setting up the stack frame, it has to traverse the list of arguments in the proper order to determine the correct offsets from the frame pointer. Both of those situations can be handled by providing a set of *cross*

Cross links.

links that connect all variables at a particular nesting level. For example, there are three nesting levels in the following fragment:

```
int Godot;
waiting( vladimir, estragon )
{
    int pozzo ;
    while( condition )
    {
        int pozzo, lucky;
    }
}
```

Godot and waiting are at the outer level, the subroutine arguments and the first pozzo comprise the second block, and the third block has the second pozzo in it. Note that the arguments are grouped with the outer local variables and that the inner pozzo shadows the outer one while the inner block is active—the two pozzos are different variables, occupying different parts of the stack frame. Figure 6.10 shows how the symbol table looks when the innermost block is being processed. I'm assuming that vladimir hashes to the same value as pozzo. Note that the subroutine name [waiting()] is considered part of the outer block, but the arguments are part of the inner block. The hash table itself is at the left—the solid arrows are links used to resolve collisions within the table. Here, the array across the top holds the heads of the cross-link chains. This array could be eliminated by passing the head-of-chain pointers as attributes on the value stack, however. You can visit all symbol-table entries for variables at a given scoping level by traversing the cross links (the dashed lines). If the head-of-chain array was a stack, you could delete the nodes for a particular level by popping the head-of-list pointer and traversing the list, deleting nodes. The current C-compiler's symbol table is organized with cross links, as just described. A hash-table element is pictured in Figure 6.10.

The top two nodes (`next` and `prev`) are maintained by the hash-table functions described in depth in Appendix A. They point at the previous and next node in the collision chain. The double indirection lets you delete an arbitrary element from the table without having to traverse the collision chain to find the predecessor. The bottom fields are managed by the maintenance-layer, symbol-table functions and are declared as the `symbol` structure in Listing 6.22 along with the symbol table itself.

The `name` field in the `symbol` structure is the name as it appears in the input. The `rname` field is the symbol's name as it appears in the output. In the case of a global variable, `rname` holds the input name with a leading underscore appended. Local static variables are given arbitrary names by the compiler because local statics in two subroutines could have the same name. The names would conflict if both were used in the output, so an arbitrary name is used. Compiler-supplied names don't have leading underscores so can't conflict with user-supplied names. The `rname` field for automatic variables and subroutine arguments is a string, which when used in an operand, evaluates to the contents of the variable in question. For example, if an `int` variable is at offset -8 from the frame pointer, the `rname` will be "fp-8".

The `level` field holds the declaration level for the symbol (0 for global symbols, 1 if they're at the outermost block in a function, and so forth). `level` is used primarily for error checking. It helps detect a duplicate declaration.

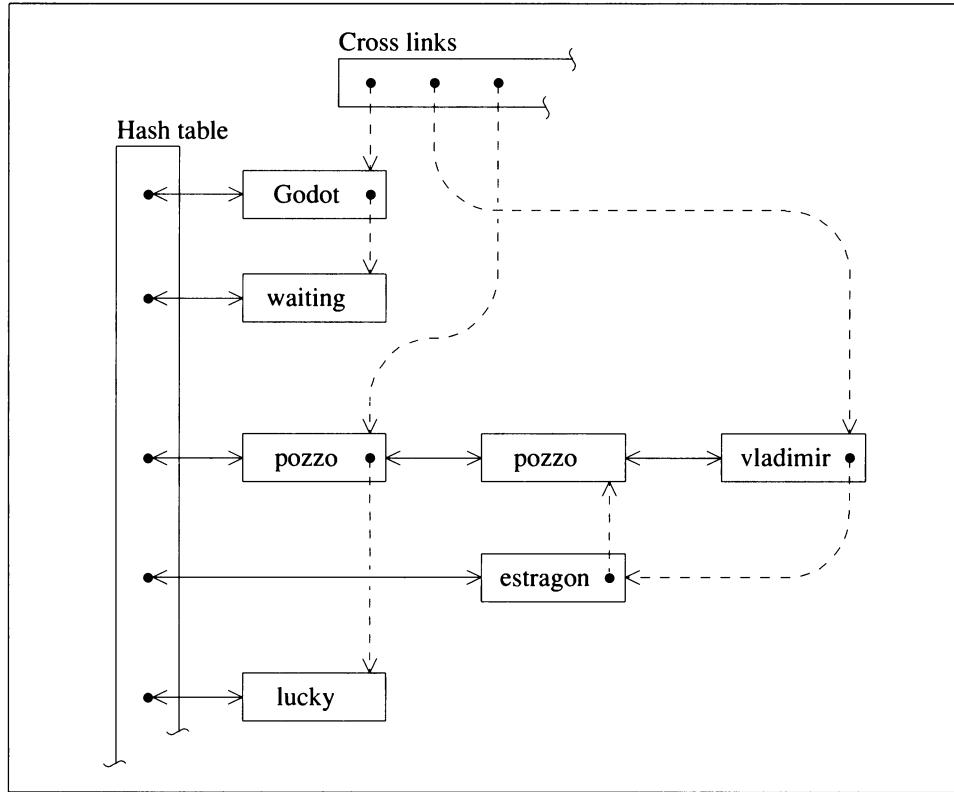
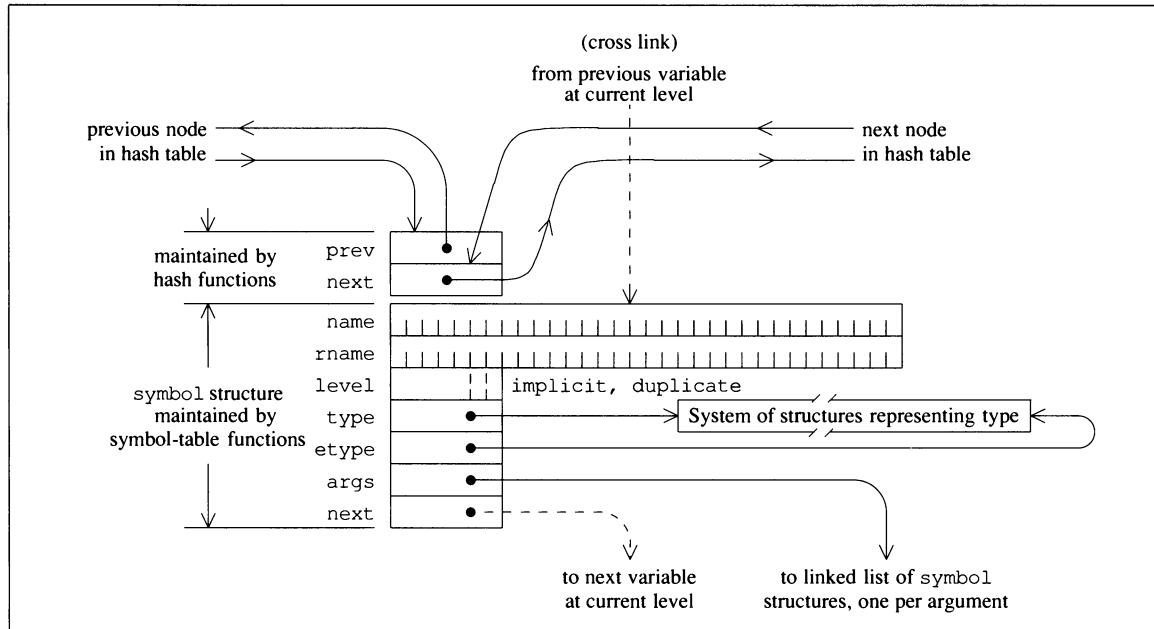
The `duplicate` bit marks the extra symbols created by duplicate declarations like the following:

`symbol.structure.`
`symbol.name,`
`symbol.rname`

Conflicting local-static
names.

`symbol.level`

`symbol.duplicate`

Figure 6.10. A Cross-linked Symbol Table**Figure 6.11.** A Symbol-Table Element

Listing 6.22. *syntab.h*—The Symbol Table

```

1  /* SYMTAB.H:  Symbol-table definitions. Note that <tools/debug.h> and
2   *               <tools/hash.h> must be #included (in that order) before the
3   *               #include for the current file.
4  */
5
6  #ifdef ALLOC                                /* Allocate variables if ALLOC defined. */
7  #  define ALLOC_CLS /* empty */
8  #else
9  #  define ALLOC_CLS  extern
10 #endif
11
12 #define NAME_MAX 32                         /* Maximum identifier length.      */
13 #define LABEL_MAX 32                         /* Maximum output-label length.   */
14
15 typedef struct symbol                      /* Symbol-table entry.           */
16 {
17     unsigned char    name [NAME_MAX+1];        /* Input variable name.          */
18     unsigned char    rname [NAME_MAX+1];        /* Actual variable name.         */
19
20     unsigned         level      : 13 ;        /* Declaration lev., field offset.*/
21     unsigned         implicit   : 1 ;        /* Declaration created implicitly.*/
22     unsigned         duplicate  : 1 ;        /* Duplicate declaration.       */
23
24     struct          link      *type;        /* First link in declarator chain.*/
25     struct          link      *etype;        /* Last link in declarator chain.*/
26     struct          symbol   *args;        /* If a funct decl, the arg list. */
27                                         /* If a var, the initializer.    */
28     struct          symbol   *next;        /* Cross link to next variable at */
29                                         /* current nesting level.       */
30 } symbol;
31
32 ALLOC_CLS HASH_TAB *Symbol_tab;           /* The actual table. */

```

```

extern int x;
...
int x = 5;

```

It's convenient not to discard these symbols as soon as the duplication is found. The duplicate symbol is not put into the symbol table, however, and this bit lets you determine that the symbol has not been inserted into the table so you can delete it at a later time.

`symbol.implicit`

The `implicit` field is used to distinguish undeclared variables from implicit function declarations created when a subroutine is used before being defined. An implicit declaration of type `int` is created when an undeclared identifier is encountered by the compiler. If that identifier is subsequently used in a function call, its type is modified to “function returning `int`” and the `implicit` bit is cleared. If a symbol is still marked `implicit` when the compiler is done with the current block, then it is an undeclared variable and an appropriate error message is printed.

`symbol.type,`
`symbol.etype,`
`symbol.args.`

`symbol.next`

The `type` and `etype` fields in the `symbol` structure point to yet another data structure (discussed shortly) that describes the object's type. `args` keeps track of function arguments until they are added to the symbol table. It is the head of a linked list of `symbol` structures, one for each argument.

Finally, the `next` pointer is the cross link to the next variable at the same nesting level. Note that this field points at the `symbol` component of the combined structure, not at the header maintained by the hash functions.

6.3.4 Representing Types—Theory

The next issue is the representation of a variable's type. If a language is simple enough, types can be represented with a simple numeric coding in the symbol structure. For example, if a language had only two types, integer and floating point, you could define two constants like this:

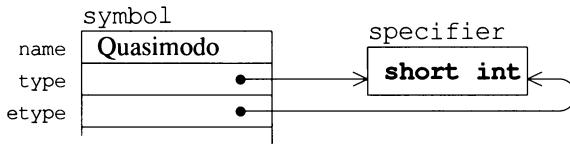
```
#define INTEGER 0
#define FLOAT 1
```

and add a field to the symbol structure that would be set to one or the other of these values. Pointers could be represented in a similar way, with a second variable keeping track of the levels of indirection. For example, given a declaration like `int ***p`, this second variable would hold 3. This sort of typing system is called a *constrained* system because there are only a limited number of possible types.

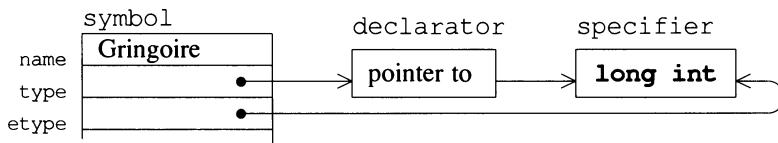
The situation is more complex in a language like C, which has an *unconstrained* typing system that permits virtually unlimited complexity in a variable declaration. A C variable's type must be represented by a system of data structures working in concert. You can see what's required by looking at how a C variable declaration is organized. Variable declarations have two parts: a *specifier* part which is a list of various keywords (`int`, `long`, `extern`, `struct`, and so forth) and a *declarator* part that is made up of the variable's name and an arbitrary number of stars, array-size specifiers (like `[10]`) and parentheses (used both for grouping and to indicate a function).

The specifier is constrained—there are only a limited number of legal combinations of keywords that can be used here—so it can be represented by a single structure. The declarator is not constrained, however—any number of stars, brackets, or parentheses are permitted in any combination. Because of this organization, a type can be represented using two kinds of structures, one representing the specifier and another representing the declarator. The type definition is a linked list of these structures, and the type field in the symbol structure that we looked at earlier points at the head of the list (etype points at the end). All type representations have exactly one specifier structure, though there can be any number of declarators (including none), and it's convenient for the specifier to be at the end of the linked list—I'll explain why in a moment.

Let's look at some examples. A simple declaration of the form `short Quasimodo;` is represented like this:



A pointer type, like `long *Gringoire;` adds a declarator structure, like this:



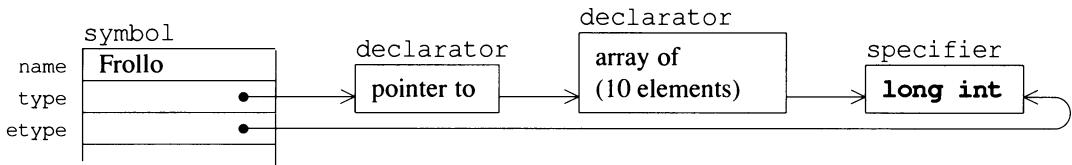
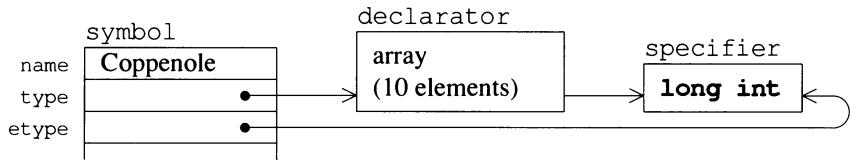
You read down the list just like you would parse the declaration in English. Gringoire is a pointer to a `long`. The symbol node holds the name, the second holds an indicator that the variable is a pointer, and the third node holds the `long`. An array of longs declared with `long Coppenole[10];` is represented as follows:

Constrained types,
numeric coding.

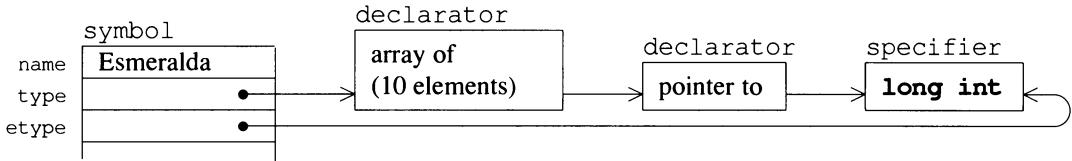
Unconstrained types.

Specifier.
Declarator.

Example type representations.



An array of pointers to longs such as **long** *Esmerelda[10] has the “pointer to” and “array of” nodes transposed, like this:



Derived types.

This system has an important characteristic that will be quite useful for generating code. The array-bracket, indirection, and structure-access operators generate temporary variables, just like the simple addition and multiplication operators discussed in earlier chapters. The types of these temporaries can be *derived* from the original type representation, often just by removing the first element in the chain. For example, in the earlier array-of-pointers-to-**long** example, an expression like Esmerelda[1] generates a temporary of type “pointer to long” that holds the referenced array element. The temporary’s type is just the original chain, less the leading “array of” node. The address-of operator (&) is handled in much the same way by adding a “pointer to” node to the left of the chain.

6.3.5 Representing Types—Implementation

Implementing declarators.

The foregoing system of declarators and specifiers can be implemented with a system of several structures. A declarator is represented with the structure in Listing 6.23. There are two fields: `dcl_type` identifies the declarator as a pointer, array, or function, using the values defined on lines 33 to 35. If the declarator is an array, `num_ele` holds the number of elements. You’ll need this number to increment a pointer to an array correctly, and to figure an array’s size.

Implementing Specifiers.

Nouns.

Adjectives.

Specifiers are represented by the structure and macros in Listing 6.24. The situation is simplified here because the current C compiler ignores the `const` and `volatile` keywords. A C specifier consists of a basic type (call it a *noun*), and various modifiers (call them *adjectives*). C supports several nouns: **char**, **int**, **float**, **double**, **struct**, **union**, and **enum**. Only four of these are necessary in the current compiler, however: **char**, **int**, **void**, and **struct**. You would add **float** to this list if floating point were supported. A **double** could be treated as a **long float**. Enumerated types are treated as if they were **int**s. Unions are treated identically to structures—all the offsets to the fields are zero, however. An **int** is implied if a noun is missing from a declaration, as in **long** **x**. C also supports one implicit identifier that is declared just by using it: a label. The `LABEL` on line 46 takes care of this sort of identifier. The basic nouns are

Listing 6.23. *syntab.h*—A Declarator

```

33 #define POINTER          0           /* Values for declarator.type. */
34 #define ARRAY            1
35 #define FUNCTION         2
36
37 typedef struct declarator
38 {
39     int dcl_type;                /* POINTER, ARRAY, or FUNCTION */
40     int num_ele;               /* If class==ARRAY, # of elements */
41 } declarator;

```

represented by the **noun** field on line 64 of Listing 6.24, and the possible values for this field are defined on lines 42 to 46.

Next come the adjectives. These fall naturally into several categories that have mutually exclusive values. The most complex of these is the storage class represented by the **sclass** field on line 65 of Listing 6.24, which can hold one of the values defined on lines 47 to 52. These values are mutually exclusive—you can't have an **extern register**, for example. Note that the **TYPEDEF** class is used only to pass information around during the declaration. **typedefs** are also marked by a bit in the first element of the type chain—it's convenient to have this information duplicated while the declaration is being processed—we'll see why in a moment. Note that **AUTO** designates anything that can be on the run-time stack, and **FIXED** is anything at a fixed address in memory, whether or not it's declared **static**. The **static** keyword is used by the compiler to decide whether or not to make the variable public when it generates the code to allocate space for that variable.

The **CONSTANT** storage class is used for two purposes. First, when you declare an **enum** type like this:

```

enum rabbits
{
    FLOPSY, MOPSEY, PETER, COTTONTAIL;
};

```

the compiler puts **int** entries for each of the elements of the enumerator list (FLOPSY, MOPSEY, PETER, and COTTONTAIL) into the symbol table and sets the **CONSTANT** attribute in the associated specifier. The **const_val** union on lines 72 to 82 of Listing 6.25-1 holds the value associated with the constant. A definition for FLOPSY is in Figure 6.12. The **v_int** field holds the numeric values of integer constants, but it's also used for string constants. When the compiler sees a string constant, it outputs a definition of the form

```
char S1[] = "contents of string" ;
```

and the numeric component of the label is stored in **v_int**—each string label has a unique numeric component. **v_struct** is used only if the current specifier describes a structure, in which case the **noun** field is set to **STRUCTURE** and **v_struct** points at yet another data structure (discussed momentarily) that describes the structure.

The **oclass** field on line 66 of Listing 6.24 remembers the C-code storage class that is actually output with the variable definition for a global variable. The contents of this field are undefined if the **sclass** field is not set to **FIXED**, and the possible values are defined on lines 54 to 59.

The **_long** field on line 67 selects either of two lengths for an integer variable. In the current application, the **long** keyword is ignored if the variable is of type **char**, and the **short** keyword is always ignored. All of this is summarized in Table 6.11. The **_long** and **_short**.

Listing 6.24. *syntab.h*—A Specifier

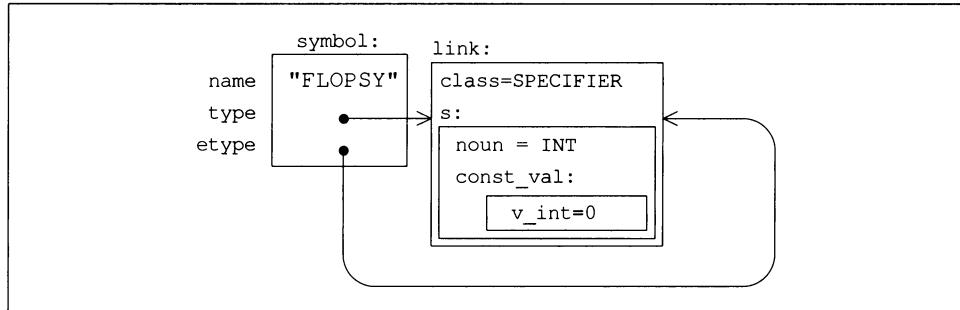
```

42 #define INT      0          /* specifier.noun. INT has the value 0 so */
43 #define CHAR     1          /* that an uninitialized structure defaults */
44 #define VOID     2          /* to int, same goes for EXTERN, below. */
45 #define STRUCTURE 3
46 #define LABEL    4
47                                         /* specifier.sclass */
48 #define FIXED    0          /* At a fixed address. */
49 #define REGISTER 1          /* In a register. */
50 #define AUTO     2          /* On the run-time stack. */
51 #define TYPEDEF   3          /* Typedef. */
52 #define CONSTANT 4          /* This is a constant. */
53
54                                         /* Output (C-code) storage class */
55 #define NO_OCLASS 0          /* No output class (var is auto). */
56 #define PUB       1          /* public */
57 #define PRI       2          /* private */
58 #define EXT       3          /* extern */
59 #define COM       4          /* common */
60
61
62 typedef struct specifier
63 {
64     unsigned noun      :3;    /* CHAR INT STRUCTURE LABEL */
65     unsigned sclass    :3;    /* REGISTER AUTO FIXED CONSTANT TYPEDEF */
66     unsigned oclass    :3;    /* Output storage class: PUB PRI COM EXT. */
67     unsigned _long     :1;    /* 1=long.      0=short. */
68     unsigned _unsigned  :1;    /* 1=unsigned.  0=signed. */
69     unsigned _static   :1;    /* 1=static keyword found in declarations. */
70     unsigned _extern   :1;    /* 1=extern keyword found in declarations. */
71
72     union
73     {
74         /* Value if constant: */
75         int           v_int;   /* Int & char values. If a string const., */
76                               /* is numeric component of the label. */
77         unsigned int   v_uint;   /* Unsigned int constant value. */
78         long          v_long;   /* Signed long constant value. */
79         unsigned long v_ulong;  /* Unsigned long constant value. */
80
81         struct structdef *v_struct; /* If this is a struct, points at a */
82                               /* structure-table element. */
83     } const_val;
84 } specifier;

```

`_unsigned` bit on line 68 is used in much the same way as `_long`. The `_extern` and `_static` bits on the next two lines remember when the equivalent keyword is found in the input as the specifier list is parsed. They are needed by the compiler to figure the output storage class after the entire specifier has been processed.

There's one final problem. A declaration list can be made up of two types of structures: it can have zero or more declarators in it, and it always has exactly one specifier. You need some way to determine if a list element is a declarator or a specifier when all you have to work with is a pointer to an element. The problem is solved by *encapsulating* the two structures into a third structure that can be either a declarator or specifier. This is done with the link structure in Listing 6.25. The `class` field tells us what the following union contains. It's either a DECLARATOR or a

Figure 6.12. An Enumerator-List Element.**Table 6.11.** Processing **long** and **short**.

input	noun	_long	length	notes
long int	INT	true	32 bits	
short int	INT	false	16 bits	same as int
int	INT	false	16 bits	
long char	CHAR	true	8 bits	same as char
short char	CHAR	false	8 bits	same as char
char	CHAR	false	8 bits	

SPECIFIER, as defined on lines 85 and 86. The next field points at the next element in the type chain. It's NULL if this is the specifier, which must come last in the chain. Finally, the tdef field is used when processing typedefs. It's used to distinguish whether a type chain was created by a **typedef** or by a normal declaration. You could treat **typedef** as a storage class, and mark it as such in the sclass field of the specifier, but it's convenient not to have to chase down the length of the chain to the specifier to get this information.

Listing 6.25. *syntab.h*—A link in the Declaration Chain

```

85 #define DECLARATOR      0
86 #define SPECIFIER       1
87
88 typedef struct link
89 {
90     unsigned class     : 1;           /* DECLARATOR or SPECIFIER */
91     unsigned tdef      : 1;           /* For typedefs. If set, current link */
92                               /* chain was created by a typedef. */
93     union
94     {
95         specifier    s;           /* If class == DECLARATOR */
96         declarator   d;           /* If class == SPECIFIER */
97     }
98     select ;
99     struct link *next;          /* Next element of chain. */
100
101 } link;
102
103 /*****
104 * Use the following p->XXX where p is a pointer to a link structure.
105 */
106

```

Listing 6.25. continued...

```

107 #define NOUN           select.s.noun
108 #define SCLASS         select.s.sclass
109 #define LONG          select.s._long
110 #define UNSIGNED      select.s._unsigned
111 #define EXTERN        select.s._extern
112 #define STATIC         select.s._static
113 #define OCLASS         select.s.oclass
114
115 #define DCL_TYPE       select.d.dcl_type
116 #define NUM_ELE         select.d.num_ele
117
118 #define VALUE          select.s.const_val
119 #define V_INT           VALUE.v_int
120 #define V_UINT          VALUE.v_uint
121 #define V_LONG          VALUE.v_long
122 #define V ULONG         VALUE.v_ulong
123 #define V_STRUCT        VALUE.v_struct
124
125 /* -----
126 * Use the following XXX(p) where p is a pointer to a link structure.
127 */
128
129 #define IS_SPECIFIER(p)  ( (p)->class == SPECIFIER )
130 #define IS_DECLARATOR(p) ( (p)->class == DECLARATOR )
131 #define IS_ARRAY(p)      ( (p)->class == DECLARATOR && (p)->DCL_TYPE==ARRAY )
132 #define IS_POINTER(p)   ( (p)->class == DECLARATOR && (p)->DCL_TYPE==POINTER )
133 #define IS_FUNCT(p)     ( (p)->class == DECLARATOR && (p)->DCL_TYPE==FUNCTION )
134 #define IS_STRUCT(p)    ( (p)->class == SPECIFIER && (p)->NOUN == STRUCTURE )
135 #define IS_LABEL(p)     ( (p)->class == SPECIFIER && (p)->NOUN == LABEL )
136
137 #define IS_CHAR(p)      ( (p)->class == SPECIFIER && (p)->NOUN == CHAR )
138 #define IS_INT(p)       ( (p)->class == SPECIFIER && (p)->NOUN == INT )
139 #define IS_UINT(p)      ( IS_INT(p) && (p)->UNSIGNED )
140 #define IS_LONG(p)      ( IS_INT(p) && (p)->LONG )
141 #define IS ULONG(p)    ( IS_INT(p) && (p)->LONG && (p)->UNSIGNED )
142 #define IS_UNSIGNED(p)  ( (p)->UNSIGNED )
143
144
145 #define IS_AGGREGATE(p) ( IS_ARRAY(p) || IS_STRUCT(p) )
146 #define IS_PTR_TYPE(p)  ( IS_ARRAY(p) || IS_POINTER(p) )
147
148 #define IS_CONSTANT(p)   (IS_SPECIFIER(p) && (p)->SCLASS == CONSTANT )
149 #define IS_TYPEDEF(p)   (IS_SPECIFIER(p) && (p)->SCLASS == TYPEDEF )
150 #define IS_INT_CONSTANT(p) (IS_CONSTANT(p) && (p)->NOUN == INT )

```

The macros on lines 107 to 123 of Listing 6.25 clean up the code a little by getting rid of some of the dots and field names. For example, if *p* is a pointer to a link, you can say *p*->*V_INT* rather than

```
p->select.d.const_val.v_int
```

to access that field.

Structures.

The foregoing system becomes even more complex when you introduce structures and unions into the picture. You need two more data structures for this purpose. First of all, the structure definitions are organized in an auxiliary symbol table called the *structure table*, declared on line 159 of Listing 6.26, below. The *v_struct* field of a specifier that describes a structure points at the structure-table element for that

structure. The table is indexed by tag name if there is one—untagged structures are assigned arbitrary names. It contains structdef structures, defined on lines 151 to 157 of Listing 6.26. The structdef contains the tag name (tag), the nesting level at the point of declaration (level), and a pointer to a linked list of field definitions (fields), each of which is a symbol structure—one symbol for each field. The level field is here so that an error message can be printed when a duplicate declaration is found; it's recycled later to hold the offset to the field from the base address of the structure. These offsets are all zero in the case of a union; that's the only difference between a structure and a union, in fact. The symbol's next field links together the field definitions. This organization means that you must use a linear search to find a field, but it lets you have an arbitrary number of fields.

structdef

Listing 6.26. *syntab.h*— Representing Structures

```

151 typedef struct structdef
152 {
153     char          tag[NAME_MAX+1]; /* Tag part of structure definition. */
154     unsigned char level;           /* Nesting level at which struct declared.*/
155     symbol        *fields;        /* Linked list of field declarations. */
156     unsigned       size;          /* Size of the structure in bytes. */
157 } structdef;
158
159 ALLOC_CLS HASH_TAB *Struct_tab; /* The actual table. */
```

Figure 6.13 gives you an idea of how a reasonably complex declaration appears in the complete symbol-table system. I've left out irrelevant fields in the figure. The declaration that generated that table is as follows:

```

struct argotiers
{
    int             (* Clopin)(); /* Function pointer */
    double          Mathias[5];
    struct argotiers *Guillaume;
    struct pstruct { int a; } Pierre;
}
gipsy;
```

Note that isolating the structdef from the field lets you correctly process a declaration like the following:

```

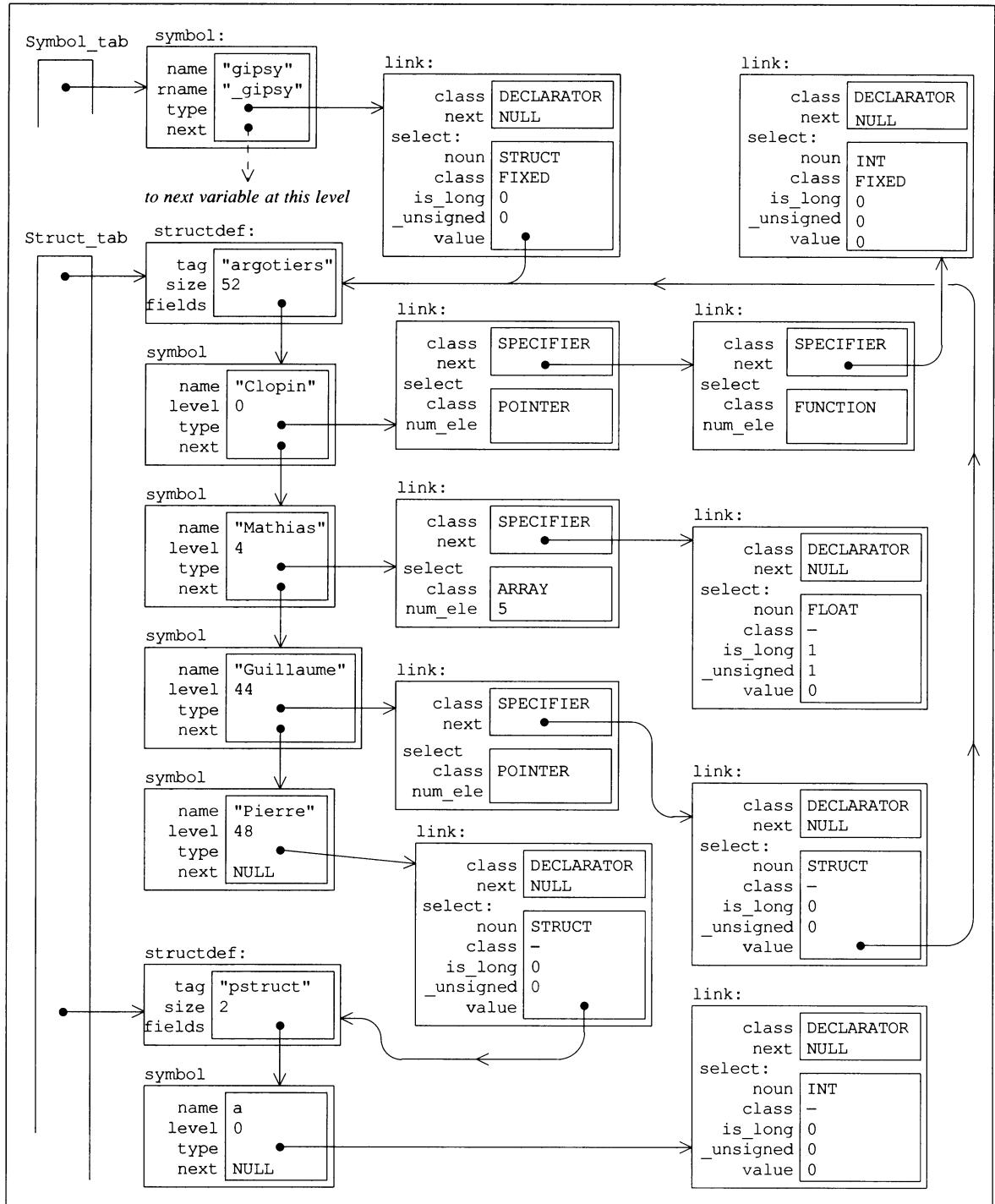
struct one { struct two *p; };
struct two { struct one *p; };
```

because you can create a struct_def for **struct** two without having to know anything about this second structure's contents. The fields can be added to the structdef when you get to the **struct** two declaration.

The final part of *syntab.h*, in which various mappings from C types to C-code types are defined, is in Listing 6.27. Note that the various WIDTH macros from *<tools/c-code.h>* are used here, so you must include *c-code.h* before including *syntab.h* in your file.

C to C-code mappings.

Figure 6.13. Representing a Structure in the Symbol Table



Listing 6.27. *syntab.h*— Sizes of Various Types

```

160 #define CSIZE    BYTE_WIDTH      /* char */
161 #define CTYPE    "byte"
162
163 #define ISIZE    WORD_WIDTH      /* int */
164 #define ITYPE    "word"
165
166 #define LSIZE    LWORD_WIDTH     /* long */
167 #define LTYPE    "lword"
168
169 #define PSIZE    PTR_WIDTH       /* pointer: 32-bit (8086 large model) */
170 #define PTYPE    "ptr"
171
172 #define STYPE    "record"        /* structure, size undefined */
173 #define ATYPE    "array"         /* array,      size undefined */

```

6.3.6 Implementing the Symbol-Table Maintenance Layer

Our next task is to build the symbol-table maintenance layer—the functions that manipulate the data structures described in the last section. The first set of functions, in Listing 6.28, take care of memory management. Three sets of similar routines are provided—they maintain the symbol, link, and structdef structure respectively. Taking the symbol-maintenance routines on lines 23 to 81 of Listing 6.28 as characteristic, the compiler creates and deletes symbols throughout the compilation process, so it's worthwhile to minimize the create-and-delete time for a node. `discard_symbol()` creates a linked list of freed nodes on line 63 rather than calling `newsym()` and `freesym()` [the hash-table-function versions of `malloc()` and `free()`] for every create and delete. `Symbol_free` points at the head of the free list; it is declared on line 15. `new_symbol()` calls `newsym()` on line 30 only if the free list is empty, otherwise it just unlinks a node from the list on line 33. The routines for the other structures work in much the same way. The only difference is that `new_link()`, since it's calling `malloc()` directly on line 96, can get nodes ten at a time to speed up the allocation process even further. `LCHUNK`—the number of nodes that are allocated at one time—is defined on line 19.

Listing 6.28. *syntab.c*— The Maintenance Layer: Memory-Management Functions

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/debug.h>
4 #include <tools/hash.h>
5 #include <tools/l.h>
6 #include <tools/compiler.h>
7 #include <tools/c-code.h>
8 #include "syntab.h"           /* Symbol-table definitions.          */
9 #include "value.h"            /* Value definitions.               */
10 #include "proto.h"           /* Prototypes for all functions in this directory. */
11 #include "label.h"            /* Labels to use for compiler-generated symbols. */
12
13 /-----*/
14
15 PRIVATE symbol *Symbol_free = NULL; /* Free-list of recycled symbols.   */
16 PRIVATE link *Link_free = NULL; /* Free-list of recycled links.     */
17 PRIVATE structdef *Struct_free = NULL; /* Free-list of recycled structdefs. */
18

```



Listing 6.28. continued...

```

19 #define LCHUNK 10           /* new_link() gets this many nodes at one shot.*/
20
21 /*-----*/
22
23 PUBLIC symbol *new_symbol( name, scope )
24 char *name;
25 int scope;
26 {
27     symbol *sym_p;
28
29     if( !Symbol_free )                                /* Free list is empty.*/
30         sym_p = (symbol *) newsym( sizeof(symbol) );
31     else
32     {
33         sym_p      = Symbol_free;
34         Symbol_free = Symbol_free->next ;
35         memset( sym_p, 0, sizeof(symbol) );
36     }
37
38     strncpy( sym_p->name, name, sizeof(sym_p->name) );
39     sym_p->level = scope;
40     return sym_p;
41 }
42
43 /*-----*/
44
45 PUBLIC void discard_symbol( sym )
46 symbol *sym;
47 {
48     /* Discard a single symbol structure and any attached links and args. Note
49      * that the args field is recycled for initializers, the process is
50      * described later on in the text (see value.c in the code), but you have to
51      * test for a different type here. Sorry about the forward reference.
52      */
53
54     if( sym )
55     {
56         if( IS_FUNCT( sym->type ) )
57             discard_symbol_chain( sym->args );          /* Function arguments. */
58         else
59             discard_value( (value *)sym->args );        /* If an initializer. */
60
61         discard_link_chain( sym->type );                /* Discard type chain. */
62
63         sym->next    = Symbol_free ;                   /* Put current symbol */
64         Symbol_free = sym;                            /* in the free list. */
65     }
66 }
67
68 /*-----*/
69
70 PUBLIC void discard_symbol_chain(sym) /* Discard an entire cross-linked */
71 symbol *sym;                         /* chain of symbols. */
72 {
73     symbol *p = sym;
74
75     while( sym )
76     {
77         p = sym->next;
78         discard_symbol( sym );

```



Listing 6.28. continued...

```
79         sym = p;
80     }
81 }
82
83 /*-----*/
84
85 PUBLIC link      *new_link( )
86 {
87     /* Return a new link. It's initialized to zeros, so it's a declarator.
88      * LCHUNK nodes are allocated from malloc() at one time.
89      */
90
91     link *p;
92     int   i;
93
94     if( !Link_free )
95     {
96         if( !(Link_free = (link *) malloc( sizeof(link) * LCHUNK ) ) )
97         {
98             yyerror("INTERNAL, new_link: Out of memory\n");
99             exit( 1 );
100        }
101
102        for( p = Link_free, i = LCHUNK; --i > 0; ++p )    /* Executes LCHUNK-1 */
103            p->next = p + 1;                                /* times. */
104
105        p->next = NULL ;
106    }
107
108    p      = Link_free;
109    Link_free = Link_free->next;
110    memset( p, 0, sizeof(link) );
111    return p;
112 }
113
114 /*-----*/
115
116 PUBLIC void      discard_link_chain( p )
117 link      *p;
118 {
119     /* Discard all links in the chain. Nothing is removed from the structure
120      * table, however. There's no point in discarding the nodes one at a time
121      * since they're already linked together, so find the first and last nodes
122      * in the input chain and link the whole list directly.
123      */
124
125     link *start ;
126
127     if( start = p )
128     {
129         while( p->next )          /* find last node */
130             p = p->next;
131
132         p->next    = Link_free;
133         Link_free = start;
134     }
135 }
136 }
```



Listing 6.28. continued...

```

137 /*-----*/
138
139 PUBLIC void discard_link( p ) /* Discard a single link. */
140 link *p;
141 {
142     p->next = Link_free;
143     Link_free = p;
144 }
145
146 /*-----*/
147
148 PUBLIC structdef *new_structdef( tag ) /* Allocate a new structdef. */
149 char *tag;
150 {
151     structdef *sdef_p;
152
153     if( !Struct_free )
154         sdef_p = (structdef *) newsym( sizeof(structdef) );
155     else
156     {
157         sdef_p = Struct_free;
158         Struct_free = (structdef *) (Struct_free->fields);
159         memset( sdef_p, 0, sizeof(structdef) );
160     }
161     strncpy( sdef_p->tag, tag, sizeof(sdef_p->tag) );
162     return sdef_p;
163 }
164
165 /*-----*/
166
167 PUBLIC void discard_structdef( sdef_p )
168 structdef *sdef_p;
169 {
170     /* Discard a structdef and any attached fields, but don't discard linked
171      * structure definitions.
172      */
173
174     if( sdef_p )
175     {
176         discard_symbol_chain( sdef_p->fields );
177
178         sdef_p->fields = (symbol *)Struct_free ;
179         Struct_free = sdef_p;
180     }
181 }
```

Subroutines in Listing 6.29 manipulate declarators: `add_declarator()` adds declarator nodes to the end of the linked list pointed to by the `type` and `etype` fields in a symbol structure. The routine is passed a symbol pointer and the declarator type (ARRAY, POINTER, or FUNCTION—declared on line 33 of Listing 6.23, page 491). Subroutines in Listing 6.30 manipulate specifiers: they copy, create, and initialize specifier links.

Listing 6.31 contains routines that manipulate entire types and linked-lists of symbols. `clone_type()`, on line 226, copies an entire type chain. This routine is used in two places. First, when a variable that uses a `typedef` rather than a standard type is declared, the type chain is copied from the symbol representing the `typedef` to the

Listing 6.29. *syntab.c*—The Maintenance Layer: Declarator Manipulation

```

182 PUBLIC void add_declarator( sym, type )
183 symbol *sym;
184 int type;
185 {
186     /* Add a declarator link to the end of the chain, the head of which is
187      * pointed to by sym->type and the tail of which is pointed to by
188      * sym->etype. *head must be NULL when the chain is empty. Both pointers
189      * are modified as necessary.
190     */
191
192     link *link_p;
193
194     if( type == FUNCTION && IS_ARRAY(sym->etype) )
195     {
196         yyerror("Array of functions is illegal, assuming function pointer\n");
197         add_declarator( sym, POINTER );
198     }
199
200     link_p = new_link();           /* The default class is DECLARATOR. */
201     link_p->DCL_TYPE = type;
202
203     if( !sym->type )
204         sym->type = sym->etype = link_p;
205     else
206     {
207         sym->etype->next = link_p;
208         sym->etype = link_p;
209     }
210 }
```

Listing 6.30. *syntab.c*—The Maintenance Layer: Specifier Manipulation

```

211 PUBLIC spec_cpy( dst, src )      /* Copy all initialized fields in src to dst.*/
212 link *dst, *src;
213 {
214
215     if( src->NOUN )          dst->NOUN      = src->NOUN      ;
216     if( src->SCLASS )        dst->SCLASS    = src->SCLASS    ;
217     if( src->LONG )          dst->LONG     = src->LONG     ;
218     if( src->UNSIGNED )      dst->UNSIGNED = src->UNSIGNED ;
219     if( src->STATIC )        dst->STATIC   = src->STATIC   ;
220     if( src->EXTERN )        dst->EXTERN   = src->EXTERN   ;
221     if( src->tdef )          dst->tdef     = src->tdef     ;
222
223     if( src->SCLASS == CONSTANT || src->NOUN == STRUCTURE)
224         memcpy( &dst->VALUE, &src->VALUE, sizeof(src->VALUE) );
225 }
```

symbol for the variable. You could keep around only one type chain in the **typedef** symbol, and make the new symbols type field point there, but this would complicate symbol deletion, because you'd have to keep track of the scope level of every link as you deleted nodes. Though this later method would be more conservative of memory, it complicates things enough so that I didn't want to use it. `clone_type()` is also used to create type chains for temporary variables, though here the copying is less defensible.

I'll discuss this second application when expression processing is discussed, below.

`the_same_type`, on line 258 of Listing 6.31, compares two type chains and returns true if they match. The storage class components of the specifier are ignored. When the third argument is true, a `POINTER` declarator is considered identical to a `ARRAY` declarator when they are found in the first positions of both type chains. This relaxation of strict type checking is necessary, again, for expression processing, because arrays are represented internally as a pointer to the first element. Strict checking is necessary when two declarations are compared for equivalence, however, so the third argument lets us disable this feature.

`get_sizeof()`, on line 302 of Listing 6.31, returns the size of an object of the type represented by its argument. Note that recursion is used on line 313 to process arrays. A declaration like this:

```
int a[10][20]
```

should return 400 (10×20 elements \times 2 bytes per element). If the current declarator is an array, the size of the current dimension is remembered, and the rest of the type chain is passed recursively to `get_sizeof()` to get the size of an array element. This could, of course, be done with a loop, but the recursion is more compact and arrays with more than three or four dimensions are rare. The other subroutines in the listing are self-explanatory.

The remainder of the maintenance layer, in Listing 6.32, is made up of subroutines that print symbols and convert fields to printable strings.

Listing 6.31. *symtab.c*—The Maintenance Layer: Type Manipulation

```

226 PUBLIC link *clone_type( tchain, endp )
227 link *tchain;           /* input: Type chain to duplicate.          */
228 link **endp;           /* output: Pointer to last node in cloned chain. */
229 {
230     /* Manufacture a clone of the type chain in the input symbol. Return a
231     * pointer to the cloned chain, NULL if there were no declarators to clone.
232     * The tdef bit in the copy is always cleared.
233 */
234
235     link *last, *head = NULL;
236
237     for( ; tchain; tchain = tchain->next )
238     {
239         if( !head )                                /* 1st node in chain. */
240             head = last = new_link();
241         else                                         /* Subsequent node. */
242         {
243             last->next = new_link();
244             last      = last->next;
245         }
246
247         memcpy( last, tchain, sizeof(*last) );
248         last->next = NULL;
249         last->tdef = 0;
250     }
251
252     *endp = last;
253     return head;
254 }
255 }
```

Listing 6.31. continued...

```

256  /*-----*/
257
258 PUBLIC int      the_same_type( p1, p2, relax )
259 link    *p1, *p2;
260 int     relax;
261 {
262     /* Return 1 if the types match, 0 if they don't. Ignore the storage class.
263      * If "relax" is true and the array declarator is the first link in the
264      * chain, then a pointer is considered equivalent to an array.
265      */
266
267     if( relax && IS_PTR_TYPE(p1) && IS_PTR_TYPE(p2) )
268     {
269         p1 = p1->next;
270         p2 = p2->next;
271     }
272
273     for(; p1 && p2 ; p1 = p1->next, p2 = p2->next)
274     {
275         if( p1->class != p2->class )
276             return 0;
277
278         if( p1->class == DECLARATOR )
279         {
280             if( (p1->DCL_TYPE != p2->DCL_TYPE) ||
281                 (p1->DCL_TYPE==ARRAY && (p1->NUM_ELE != p1->NUM_ELE)) )
282                 return 0;
283         }
284         else
285             /* this is done last */
286         {
287             if( (p1->NOUN      == p2->NOUN      ) && (p1->LONG == p2->LONG ) &&
288                 (p1->UNSIGNED == p2->UNSIGNED ) )
289             {
290                 return ( p1->NOUN==STRUCTURE ) ? p1->V_STRUCT == p2->V_STRUCT
291                               : 1 ;
292             }
293             return 0;
294         }
295     }
296     yyerror("INTERNAL the_same_type: Unknown link class\n");
297     return 0;
298 }
299
300 /*-----*/
301
302 PUBLIC int      get_sizeof( p )
303 link    *p;
304 {
305     /* Return the size in bytes of an object of the type pointed to by p.
306      * Functions are considered to be pointer sized because that's how they're
307      * represented internally.
308      */
309
310     int size;
311
312     if( p->class == DECLARATOR )
313         size = (p->DCL_TYPE==ARRAY) ? p->NUM_ELE * get_sizeof(p->next) : PSIZE;

```



Listing 6.31. continued...

```

314     else
315     {
316         switch( p->NOUN )
317         {
318             case CHAR:      size = CSIZE;                      break;
319             case INT:       size = p->LONG ? LSIZE : ISIZE; break;
320             case STRUCTURE: size = p->V_STRUCT->size;        break;
321             case VOID:      size = 0;                         break;
322             case LABEL:     size = 0;                         break;
323         }
324     }
325
326     return size;
327 }
328
329 /*-----*/
330
331 PUBLIC symbol *reverse_links( sym )
332 symbol *sym;
333 {
334     /* Go through the cross-linked chain of "symbols", reversing the direction
335      * of the cross pointers. Return a pointer to the new head of chain
336      * (formerly the end of the chain) or NULL if the chain started out empty.
337      */
338
339     symbol *previous, *current, *next;
340
341     if( !sym )
342         return NULL;
343
344     previous = sym;
345     current = sym->next;
346
347     while( current )
348     {
349         next           = current->next;
350         current->next = previous;
351         previous     = current;
352         current       = next;
353     }
354
355     sym->next = NULL;
356     return previous;
357 }
```

Listing 6.32. syntab.c—The Maintenance Layer: Symbol-Printing Functions

```

358 PUBLIC char    *sclass_str( class )    /* Return a string representing the */
359 int    class;                      /* indicated storage class.          */
360 {
361     return class==CONSTANT    ? "CON" :
362     class==REGISTER     ? "REG" :
363     class==TYPEDEF      ? "TYP" :
364     class==AUTO         ? "AUT" :
365     class==FIXED        ? "FIX" : "BAD SCLASS" ;
366 }
367 
```

Listing 6.32. continued...

```

368  /*-----*/
369
370 PUBLIC char *oclass_str( class )      /* Return a string representing the */
371 int   class;                      /* indicated output storage class. */
372 {
373     return class==PUB ? "PUB" :
374         class==PRI ? "PRI" :
375         class==COM ? "COM" :
376         class==EXT ? "EXT" : "(NO OCLS)" ;
377 }
378
379 /*-----*/
380
381 PUBLIC char *noun_str( noun )        /* Return a string representing the */
382 int   noun;                      /* indicated noun. */
383 {
384     return noun==INT      ? "int" :
385         noun==CHAR      ? "char" :
386         noun==VOID      ? "void" :
387         noun==LABEL     ? "label" :
388         noun==STRUCTURE ? "struct" : "BAD NOUN" ;
389 }
390
391 /*-----*/
392
393 PUBLIC char *attr_str( spec_p )      /* Return a string representing all */
394 specifier *spec_p;                /* attributes in a specifier other */
395 {                                /* than the noun and storage class. */
396     static char str[5];
397
398     str[0] = ( spec_p->_unsigned ) ? 'u' : '.';
399     str[1] = ( spec_p->_static ) ? 's' : '.';
400     str[2] = ( spec_p->_extern ) ? 'e' : '.';
401     str[3] = ( spec_p->_long ) ? 'l' : '.';
402     str[4] = '\0';
403
404     return str;
405 }
406
407 /*-----*/
408
409 PUBLIC char *type_str ( link_p )    /* Return a string representing the */
410 link   *link_p;                  /* type represented by the link chain. */
411 {
412     int           i;
413     static char target [ 80 ];
414     static char buf    [ 64 ];
415     int           available = sizeof(target) - 1;
416
417     *buf    = '\0';
418     *target = '\0';
419
420     if( !link_p )
421         return "(NULL)";
422
423     if( link_p->tdef )
424     {
425         strcpy( target, "tdef " );
426         available -= 5;
427     }

```



Listing 6.32. continued...

```

428
429     for( ; link_p ; link_p = link_p->next )
430     {
431         if( IS_DECLARATOR(link_p) )
432         {
433             switch( link_p->DCL_TYPE )
434             {
435                 case POINTER:    i = sprintf(buf, "*");                                break;
436                 case ARRAY:     i = sprintf(buf, "[%d]", link_p->NUM_ELE); break;
437                 case FUNCTION: i = sprintf(buf, "()");                                break;
438                 default:       i = sprintf(buf, "BAD DECL" );                      break;
439             }
440         }
441         else /* it's a specifier */
442         {
443             i = sprintf( buf, "%s %s %s %s", noun_str( link_p->NOUN ),
444                                     sclass_str( link_p->SCLASS ),
445                                     oclass_str( link_p->OCLASS ),
446                                     attr_str( &link_p->select.s));
447
448         if( link_p->NOUN==STRUCTURE || link_p->SCLASS==CONSTANT )
449         {
450             strncat( target, buf, available );
451             available -= i;
452
453             if( link_p->NOUN != STRUCTURE )
454                 continue;
455             else
456                 i = sprintf(buf, " %s", link_p->V_STRUCT->tag ?
457                             link_p->V_STRUCT->tag : "untagged");
458         }
459     }
460
461     strncat( target, buf, available );
462     available -= i;
463 }
464
465     return target;
466 }
467
468 -----*/
469
470 PUBLIC char *tconst_str( type )
471 link      *type;                                /* Return a string representing the value */
472 {                                                 /* field at the end of the specified type */
473     static char buf[80];                          /* (which must be char*, char, int, long,
474                                                 /* unsigned int, or unsigned long). Return */
475     buf[0] = '?';                                /* "?" if the type isn't any of these. */
476     buf[1] = '\0';
477
478     if( IS_POINTER(type) && IS_CHAR(type->next) )
479     {
480         sprintf( buf, "%s%d", L_STRING, type->next->V_INT );
481     }
482     else if( !(IS_AGGREGATE(type) || IS_FUNCT(type)) )
483     {
484         switch( type->NOUN )
485         {

```



Listing 6.32. continued...

```

486     case CHAR:      sprintf( buf, "'%s' (%d)", bin_to_ascii(
487                           type->UNSIGNED ? type->V_UINT
488                           : type->V_INT,1),
489                           type->UNSIGNED ? type->V_UINT
490                           : type->V_INT,1 ) );
491     break;
492
493     case INT:       if( type->LONG )
494     {
495         if( type->UNSIGNED )
496             sprintf(buf, "%luL", type->V ULONG);
497         else
498             sprintf(buf, "%ldL", type->V LONG );
499     }
500     else
501     {
502         if( type->UNSIGNED )
503             sprintf( buf, "%u", type->V_UINT);
504         else
505             sprintf( buf, "%d", type->V_INT );
506     }
507     break;
508 }
509 }
510
511 if( *buf == '?' )
512     yyerror("Internal, tconst_str: Can't make constant for type %s\n",
513                         type_str( type ));
514 return buf;
515 }
516
517 /*-----*/
518
519 PUBLIC char    *sym_chain_str( chain )
520 symbol *chain;
521 {
522     /* Return a string listing the names of all symbols in the input chain (or
523      * a constant value if the symbol is a constant). Note that this routine
524      * can't call type_str() because the second-order recursion messes up the
525      * buffers. Since the routine is used only for occasional diagnostics, it's
526      * not worth fixing this problem.
527 */
528
529     int          i;
530     static char   buf[80];
531     char         *p      = buf;
532     int          avail   = sizeof( buf ) - 1;
533
534     *buf = '\0';
535     while( chain && avail > 0 )
536     {
537         if( IS_CONSTANT(chain->etype) )
538             i = sprintf( p, "%0.*s", avail - 2, "const" );
539         else
540             i = sprintf( p, "%0.*s", avail - 2, chain->name );
541
542         p      += i;
543         avail -= i;
544

```



Listing 6.32. continued...

```

545         if( chain = chain->next )
546     {
547         *p++ = ',';
548         i -= 2;
549     }
550 }
551
552     return buf;
553 }
554
555 /*-----*/
556
557 PRIVATE void psym( sym_p, fp )           /* Print one symbol to fp. */
558 symbol *sym_p;
559 FILE *fp;
560 {
561     fprintf(fp, "%-18.18s %-18.18s %2d %p %s\n",
562             sym_p->name,
563             sym_p->type ? sym_p->rname : "-----",
564             sym_p->level,
565             (void far *)sym_p->next,
566             type_str(sym_p->type) );
567 }
568
569 /*-----*/
570
571 PRIVATE void pstruct( sdef_p, fp )        /* Print a structure definition to fp */
572 structdef *sdef_p;                      /* including all the fields & types. */
573 FILE *fp;
574 {
575     symbol *field_p;
576
577     fprintf(fp, "struct <%s> (level %d, %d bytes):\n",
578             sdef_p->tag, sdef_p->level, sdef_p->size);
579
580     for( field_p = sdef_p->fields; field_p; field_p=field_p->next )
581     {
582         fprintf(fp, "%-20s (offset %d) %s\n",
583                 field_p->name, field_p->level, type_str(field_p->type));
584     }
585 }
586
587 /*-----*/
588
589 PUBLIC print_syms( filename )           /* Print the entire symbol table to */
590 char *filename;                         /* the named file. Previous contents */
591 {                                       /* of the file (if any) are destroyed.*/
592     FILE *fp;
593
594     if( !(fp = fopen(filename,"w")) )
595         yyerror("Can't open symbol-table file\n");
596     else
597     {
598         fprintf(fp, "Attributes in type field are: upel\n" );
599         fprintf(fp, "    unsigned (. for signed)-----+|||\n" );
600         fprintf(fp, "    private  (. for public)-----+|||\n" );
601         fprintf(fp, "    extern   (. for common)-----+|\n" );
602         fprintf(fp, "    long     (. for short )-----+\\n\\n" );
603

```



Listing 6.32. continued...

```

604     fprintf(fp,"name           rname          lev   next   type\n");
605     ptab( Symbol_tab, psym, fp, 1 );
606
607     fprintf(fp, "\nStructure table:\n\n");
608     ptab( Struct_tab, pstruct, fp, 1 );
609
610     fclose( fp );
611 }
612 }
```

6.4 The Parser: Configuration

I'll start looking at the compiler proper with the configuration parts of the **occs** input file. Listing 6.33 shows most of the header portion of the **occs** input file. The **%union** directive on lines 21 to 30 of Listing 6.33 types the value stack as a C union containing the indicated fields. Some of these fields are pointers to symbol-table structures, others to structures that are discussed below, and the last two fields are simple-variable attributes—**num** is for integer numbers, **ascii** for ASCII characters. Further down in the table, directives like **<ascii>** tell **occs** which fields to use for particular symbols. For example, an **ASSIGNOP** has an **<ascii>** attribute because the **<ascii>** directive precedes its definition on line 63. The **ascii** field of the union is used automatically if **\$N** in a production corresponds to an **ASSIGNOP**. You don't have to say **\$N.ascii**.

%union, used for C compiler's value stack.

Listing 6.33. c.y— occs Input File: Definitions Section

```

1  %{
2  #include <stdio.h>
3  #include <stdlib.h>
4  #include <signal.h>
5  #include <tools/debug.h>      /* Misc. macros. (see Appendix A) */
6  #include <tools/hash.h>       /* Hash-table support. (see Appendix A) */
7  #include <tools/compiler.h>    /* Prototypes for comp.lib functions. */
8  #include <tools/l.h>          /* Prototypes for l.lib functions. */
9
10 #define ALLOC      /* Define ALLOC to create symbol table in symtab.h. */
11 #include "symtab.h" /* Definitions for the symbol-table. */
12 #include "value.h"  /* Definitions used for expression processing. */
13 #include "label.h"  /* Prefixes used for compiler-generated labels. */
14 #include "switch.h" /* Definitions used for switch processing. */
15 #include "proto.h"  /* Function prototypes for all .c files used by the */
16 %}                  /* parser. It is not printed anywhere in the book, */
17                      /* but is included on the distribution disk. */
18
19 /*-----*/
20
21 %union {
22     char      *p_char;
23     symbol    *p_sym;
24     link      *p_link;
25     structdef *p_sdef;
26     specifier *p_spec;
27     value     *p_val;
28     int       num;           /* Make short if sizeof(int) > sizeof(int*) */
29     int       ascii;
```



Listing 6.33. continued...

```

30 }
31
32 /*-----*/
33
34 %term STRING          /* String constant. */
35 %term ICON            /* Integer or long constant including '\t', etc. */
36 %term FCON            /* Floating-point constant. */
37
38 %term TYPE            /* int char long float double signed unsigned short */
39             /* const volatile void */
40 %term <ascii> STRUCT /* struct union */
41 %term ENUM             /* enum */
42
43 %term RETURN GOTO
44 %term IF ELSE
45 %term SWITCH CASE DEFAULT
46 %term BREAK CONTINUE
47 %term WHILE DO FOR
48 %term LC RC           /* { } */
49 %term SEMI             /* ; */
50 %term ELLIPSIS         /* ... */
51
52 /* The attributes used below tend to be the sensible thing. For example, the
53 * ASSIGNOP attribute is the operator component of the lexeme; most other
54 * attributes are the first character of the lexeme. Exceptions are as follows:
55 *      token      attribute
56 *      RELOP >    '>'
57 *      RELOP <    '<'
58 *      RELOP >=   'G'
59 *      RELOP <=   'L'
60 */
61
62 %left COMMA           /* , */
63 %right EQUAL <ascii> ASSIGNOP /* = /= %= += -= <<= >>= &= |= ^= */
64 %right QUEST COLON    /* ? : */
65 %left OROR             /* || */
66 %left ANDAND           /* && */
67 %left OR                /* | */
68 %left XOR               /* ^ */
69 %left AND               /* & */
70 %left <ascii> EQUOP    /* == != */
71 %left <ascii> RELOP    /* <= >= < > */
72 %left <ascii> SHIFTOP  /* >> << */
73 %left PLUS MINUS      /* + - */
74 %left STAR <ascii> DIVOP /* * / % */
75 %right SIZEOF <ascii> UNOP INCOP /* sizeof ! ~ ++ -- */
76 %left LB RB LP RP <ascii> STRUCTOP /* [ ] ( ) . -> */
77
78
79             /* These attributes are shifted by the scanner. */
80 %term <p_sym> TTYPE    /* Name of a type created with a previous typedef. */
81             /* Attribute is a pointer to the symbol table */
82             /* entry for that typedef. */
83 %nonassoc <ascii> CLASS /* extern register auto static typedef. Attribute */
84             /* is the first character of the lexeme. */
85 %nonassoc <p_sym> NAME  /* Identifier or typedef name. Attribute is NULL */
86             /* if the symbol doesn't exist, a pointer to the */

```



Listing 6.33. continued...

```

87                                     /* associated "symbol" structure, otherwise.      */
88
89 %nonassoc ELSE
90
91     /* This gives a high precedence to ELSE to suppress
92     * the shift/reduce conflict error message in:
93     *   s -> IF LP expr RP expr | IF LP expr RP s ELSE s
94     * The precedence of the first production is the same
95     * as RP. Making ELSE higher precedence forces
96     * resolution in favor of the shift.
97 */
98 %type <num> args const_expr test
99
100 %type <p_sym> ext_decl_list ext_decl def_list def_decl_list decl
101 %type <p_sym> var_decl funct_decl local_defs new_name name enumerator
102 %type <p_sym> name_list var_list param_declaration abs_decl abstract_decl
103 %type <p_val> expr binary non_comma_expr unary initializer initializer_list
104 %type <p_val> or_expr and_expr or_list and_list
105
106 %type <p_link> type specifiers opt_specifiers type_or_class typeSpecifier
107 %type <p_sdef> opt_tag tag structSpecifier
108 %type <p_char> string_const target
109
110 /*-----
111  * Global and external variables. Initialization to zero for all globals is
112  * assumed. Since occs -a and -p is used, these variables may not be private.
113  */
114
115 %}
116 #ifdef YYACTION
117
118 extern char *yytext;                                /* generated by LeX */
119 extern const int yylineno, yyleng;
120 %}

```

There are several other variables declared in this part of the input file. I'll discuss these variables, below, as they are used.

```

180 %
181 #endif
182 %

```

Lines 34 to 50 of Listing 6.33 contain definitions for those tokens that aren't used in expressions, so don't require precedence or associativity information. Lines 62 to 76 comprise a precedence chart in which the other tokens are defined. Lowest-precedence operators are at the top of the list, and tokens on the same line are at the same precedence level. Some of these tokens have the `<ascii>` attribute associated with them. These tokens all represent multiple input symbols. The attribute, created by the lexical analyzer and passed to the parser using the `yyval` mechanism described in Appendix E, serves to identify which input symbol has been scanned. You can access the attribute from within a production by using the normal \$ mechanism—if a symbol with an `<ascii>` attribute is at \$1 in the production, then that attribute can be referenced by using \$1 in an action. Most `<ascii>` attributes are the first character of the lexeme—a few exceptions are listed in the comment on line 52 of Listing 6.33.

TYPE, CLASS, NAME, and ELSE are assigned precedence levels on lines 79 to 95 to eliminate various shift/reduce conflicts inherent in the grammar. (See Appendix E for a

Tokens not in expressions. Single tokens represent multiple symbols. Attributes differentiate them.

TYPE, CLASS, NAME, and ELSE have precedence.

description of what's going on.) The NAME and TTYPE tokens also have pointer-to-symbol-table-entry attributes associated with them. (A TTYPE is a type created by a previous **typedef**.) The lexical analyzer uses the symbol table to distinguish identifiers from user-defined types. It returns a NULL attribute if the input symbol isn't in the table, otherwise the attribute is a pointer to the appropriate symbol-table entry. I'll demonstrate how this is done in a moment.

%type

The %type directives on lines 96 to 106 of Listing 6.33 attach various attributes to nonterminal names. The abbreviations that are used in the names are defined in Table 6.12. The remainder of the header holds variable definitions, most of which aren't shown in Listing 6.33. They're discussed later along with the code that actually uses them.

Table 6.12. Abbreviations Used in Nonterminal Names

Abbreviation	Meaning	Abbreviation	Meaning
abs	abstract	expr	expression
arg	argument	ext	external
const	constant	opt	optional
decl	declarator	param	parameter
def	definition	struct	structure

SIGINT handling, clean-up.

Create temporary-files for code, data, and bss segments.

Skipping past the productions themselves for a moment, the end of the **occs** input file starts in Listing 6.34. This section contains various initialization and debugging routines that are used by the parser. The **occs** parser's initialization function starts on line 1306 of Listing 6.34. It creates temporary files and installs a signal handler that deletes the temporaries if SIGINT. is issued while the compiler is running. (SIGINT is issued when a Ctrl-C or Ctrl-Break is encountered under MS-DOS. Some UNIX systems also recognize DEL or Rubout.) **init_output_streams()**, starting on line 1251, creates three temporary files to hold output for the code, data, and bss segments. These files make it easier to keep the various segments straight as code is generated. The cleanup code that starts on line 1323 merges the three temporaries together at the end of the compilation process by renaming the data file and then appending the other two files.

Listing 6.34. *c.y*—**occs** Input File: Temporary-File Creation

```

1240 %%  

1241 extern char *mktemp();  

1242 extern FILE *yycodeout, *yydataout, *yybssout; /* In occs output file. */  

1243 extern yyprompt(); /* In debugger. */  

1244  

1245 #define OFILE_NAME "output.c" /* Output file name. */  

1246 char *Bss; /* Name of BSS temporary file. */  

1247 char *Code; /* Name of Code temporary file. */  

1248 char *Data; /* Name of Data temporary file. */  

1249 /*-----*/  

1250  

1251 PRIVATE void init_output_streams( p_code, p_data, p_bss)  

1252 char **p_code, **p_data, **p_bss;  

1253 {  

1254     /* Initialize the output streams, making temporary files as necessary.  

1255      * Note that the ANSI tmpfile() or the UNIX mkstemp() functions are both  

1256      * better choices than the mktemp()/fopen() used here because another  

1257      * process could, at least in theory, sneak in between the two calls.  

1258      * Since mktemp uses the process id as part of the file name, this  

1259      * is not much of a problem, and the current method is more portable  

1260      * than tmpfile() or mkstemp(). Be careful in a network environment. */
1261 }
```



Listing 6.34. continued...

```

1261      */
1262
1263  if( !(*p_code = mktemp("ccXXXXXX")) || !(*p_data = mktemp("cdXXXXXX")) || 
1264      !(*p_bss  = mktemp("cbXXXXXX"))
1265    )
1266  {
1267      yyerror("Can't create temporary-file names");
1268      exit( 1 );
1269  }
1270
1271  if( ! (yycodeout=fopen(*p_code, "w")) || !(yydataout=fopen(*p_data, "w")) || 
1272      !(ybbssout =fopen( *p_bss, "w"))
1273    )
1274  {
1275      perror("Can't open temporary files");
1276      exit( 1 );
1277  }
1278 }
1279
1280 /*-----*/
1281 void     (*Osig)(); /* Previous SIGINT handler. Initialized in yy_init_occs().*/
1282
1283 PRIVATE sigint_handler()
1284 {
1285     /* Ctrl-C handler. Note that the debugger raises SIGINT on a 'q' command,
1286      * so this routine is executed when you exit the debugger. Also, the
1287      * debugger's own SIGINT handler, which cleans up windows and so forth, is
1288      * installed before yy_init_occs() is called. It's called here to clean up
1289      * the environment if necessary. If the debugger isn't installed, the call
1290      * is harmless.
1291  */
1292
1293     signal( SIGINT, SIG_IGN );
1294     clean_up( );
1295     unlink( OFILE_NAME );
1296     (*Osig)();
1297     exit(1);           /* Needed only if old signal handler returns. */
1298 }
1299
1300 /*-----*/
1301 sym_cmp( s1, s2 ) symbol    *s1, *s2; { return strcmp( s1->name, s2->name );}
1302 struct_cmp( s1, s2 ) structdef *s1, *s2; { return strcmp( s1->tag, s2->tag );}
1303 sym_hash( s1 ) symbol    *s1;       { return hash_add(s1->name ); }
1304 struct_hash( s1 ) structdef *s1;     { return hash_add(s1->tag ) ; }
1305
1306 PUBLIC void     yy_init_occs( val )
1307 yystype *val;
1308 {
1309     yycomment("Initializing\n");
1310
1311     Osig = signal( SIGINT, SIG_IGN );
1312     init_output_streams( &Code, &Data, &Bss );
1313     signal( SIGINT, sigint_handler );
1314
1315     val->p_char = "----";          /* Provide attribute for the start symbol. */
1316
1317     Symbol_tab = maketab( 257, sym_hash,      sym_cmp      );
1318     Struct_tab = maketab( 127, struct_hash,   struct_cmp );
1319 }
1320

```



Listing 6.34. continued...

```

1321 /*-----*/
1322
1323 PRIVATE void clean_up()
1324 {
1325     /* Cleanup actions. Mark the ends of the various segments, then merge the
1326     * three temporary files used for the code, data, and bss segments into a
1327     * single file called output.c. Delete the temporaries. Since some compilers
1328     * don't delete an existing file with a rename(), it's best to assume the
1329     * worst. It can't hurt to delete a nonexistent file, you'll just get an
1330     * error back from the operating system.
1331 */
1332
1333     extern FILE *yycodeout, *yydataout, *yybssout;
1334
1335     signal ( SIGINT, SIG_IGN );
1336     fclose ( yycodeout );
1337     fclose ( yydataout );
1338     fclose ( yybssout );
1339     unlink ( OFILE_NAME ); /* delete old output file (ignore EEXIST) */
1340
1341     if( rename( Data, OFILE_NAME ) )
1342         yyerror("Can't rename temporary (%s) to %s\n", Data, OFILE_NAME );
1343     else
1344     {
1345         movefile( OFILE_NAME, Bss , "a" ); /* files to the end of the */
1346         movefile( OFILE_NAME, Code, "a" ); /* output file and delete the */
1347     } /* temporary files. movefile() */
1348 } /* is in appendix A. */

```

Print the value stack,
yypstk().

The remainder of the **occs** input file is in Listing 6.35. **yypstk()**, starting on line 1423, is used by the **occs** interactive debugger to print the value-stack contents. The routine is passed a pointer to the value stack item and a pointer to the equivalent string on the symbolic, debugging stack (the middle column in the stack window). It uses this string to identify which of the fields in the union are active, by searching the table on lines 1357 to 1415 with the **bsearch()** call on line 1432. The following switch then prints out stack item.

main(), yyhook_a(),
yyhook_b()

Listing 6.36 finishes up the start-up code with a **main()** function. The **yyhook_a()** routine is a debugger hook that prints the symbol table when a Ctrl-A is issued at the IDE's command prompt. The **yyhook_b()** lets you enable or disable the run-time trace feature discussed along with the **gen()** subroutine, below, from the debugger.

Listing 6.35. *c.y*—*occs* Input File: Initialization and Cleanup

```

1349 enum union_fields { NONE, P_SYM, P_LINK, P_SDEF, P_FIELD, P_SPEC,
1350                                P_CHAR, P_VALUE, SYM_CHAIN, ASCII, NUM } ;
1351 typedef struct tabtype
1352 {
1353     char *sym;
1354     enum union_fields case_val;
1355 } tabtype;
1356
1357 tabtype Tab[] =
1358 {
1359     /* nonterminal           field      */
1360     /*      name            in %union   */
1361
1362     { "ASSIGNOP",          ASCII      },
1363     { "CLASS",             ASCII      },
1364     { "DIVOP",              ASCII      },
1365     { "EQUOP",              ASCII      },
1366     { "INCOP",              ASCII      },
1367     { "NAME",               P_SYM      },
1368     { "RELOP",              ASCII      },
1369     { "SHIFTOP",             ASCII      },
1370     { "STRUCT",              ASCII      },
1371     { "STRUCTOP",             ASCII      },
1372     { "TTYPE",               P_SYM      },
1373     { "UNOP",                ASCII      },
1374     { "abs_decl",             P_SYM      },
1375     { "abstract_decl",        P_SYM      },
1376     { "and_expr",              P_VALUE    },
1377     { "and_list",              P_VALUE    },
1378     { "args",                  NUM        },
1379     { "binary",                P_VALUE    },
1380     { "const_expr",              NUM        },
1381     { "decl",                  P_SYM      },
1382     { "decl_list",             SYM_CHAIN },
1383     { "def",                   SYM_CHAIN },
1384     { "def_list",              SYM_CHAIN },
1385     { "enumerator",             P_SYM      },
1386     { "expr",                  P_VALUE    },
1387     { "ext_decl",              P_SYM      },
1388     { "ext_decl_list",         SYM_CHAIN },
1389     { "funct_decl",             P_SYM      },
1390     { "initializer",             P_VALUE    },
1391     { "initializer_list",        P_VALUE    },
1392     { "local_defs",             SYM_CHAIN },
1393     { "name",                  P_SYM      },
1394     { "name_list",              SYM_CHAIN },
1395     { "new_name",                P_SYM      },
1396     { "non_comma_expr",         P_VALUE    },
1397     { "opt_specifiers",        P_LINK      },
1398     { "opt_tag",                 P_SDEF      },
1399     { "or_expr",                  P_VALUE    },
1400     { "or_list",                  P_VALUE    },
1401     { "param_declaration",       SYM_CHAIN },
1402     { "specifiers",              P_LINK      },
1403     { "string_const",             P_CHAR      },
1404     { "structSpecifier",        P_SDEF      },
1405     { "tag",                     P_SDEF      },
1406     { "test",                      NUM        },
1407     { "type",                     P_LINK      }

```



Listing 6.35. continued...

```

1408         { "type_or_class",           P_LINK      },
1409         { "typeSpecifier",          P_LINK      },
1410         { "unary",                 P_VALUE     },
1411         { "var_decl",              P_SYM       },
1412         { "var_list",              SYM_CHAIN   },
1413         { "{72}",                  NUM        },
1414         { "{73}",                  P_VALUE     }
1415     };
1416
1417 tcmp( p1, p2 )
1418 tabtype *p1, *p2;
1419 {
1420     return( strcmp(p1->sym, p2->sym) );
1421 }
1422
1423 char    *yypstk( v, name )
1424 YYSTYPE *v;                      /* Pointer to value-stack item. */
1425 char    *name;                  /* Pointer to debug-stack item. */
1426 {
1427     static char buf[128];
1428     char        *text;
1429     tabtype     *tp, template;
1430
1431     template.sym = name;
1432     tp = (tabtype *) bsearch( &template, Tab, sizeof(Tab)/sizeof(*Tab),
1433                               sizeof(*Tab), tcmp);
1434
1435     sprintf( buf, "%04x ", v->num );      /* The first four characters in the */
1436     text = buf + 5;                         /* string are the numeric value of */
1437                                         /* the current stack element. */
1438                                         /* Other text is written at "text". */
1439     switch( tp ? tp->case_val : NONE )
1440     {
1441     case SYM_CHAIN:
1442         sprintf( text, "sym chain: %s",
1443                   v->p_sym ? sym_chain_str(v->p_sym) : "NULL" );
1444         break;
1445     case P_SYM:
1446         if( ! v->p_sym )
1447             sprintf( text, "symbol: NULL" );
1448
1449         else if( IS_FUNCT(v->p_sym->type) )
1450             sprintf( text, "symbol: %s(%s)=%s %1.40s",
1451                       v->p_sym->name,
1452                       sym_chain_str( v->p_sym->args ),
1453                       v->p_sym->type && *(v->p_sym->rname) ?
1454                                         v->p_sym->rname : "",
1455                                         type_str(v->p_sym->type) );
1456         else
1457             sprintf( text, "symbol: %s=%s %1.40s",
1458                       v->p_sym->name,
1459                       v->p_sym->type && *(v->p_sym->rname) ?
1460                                         v->p_sym->rname : "",
1461                                         type_str(v->p_sym->type) );
1462         break;
1463     case P_SPEC:
1464         if( !v->p_spec )
1465             sprintf( text, "specifier: NULL" );

```



Listing 6.35. continued...

```

1466           else
1467               sprintf( text, "specifier: %s %s", attr_str( v->p_spec ),
1468                               noun_str( v->p_spec->noun ) );
1469           break;
1470     case P_LINK:
1471         if( !v->p_link )
1472             sprintf( text, "specifier: NULL" );
1473         else
1474             sprintf( text, "link: %1.50s", type_str(v->p_link) );
1475         break;
1476     case P_VALUE:
1477         if( !v->p_val )
1478             sprintf( text, "_value: NULL" );
1479         else
1480         {
1481             sprintf( text, "%cvalue: %s %c/%u %1.40s",
1482                     v->p_val->lvalue ? 'l' : 'r' ,
1483                     * (v->p_val->name) ? v->p_val->name : "--",
1484                     v->p_val->is_tmp ? 't' : 'v' ,
1485                     v->p_val->offset,
1486                     type_str( v->p_val->type ) );
1487         }
1488         break;
1489     case P_SDEF:
1490         if( !v->p_sdef )
1491             sprintf( text, "structdef: NULL" );
1492         else
1493             sprintf( text, "structdef: %s lev %d, size %d",
1494                     v->p_sdef->tag,
1495                     v->p_sdef->level,
1496                     v->p_sdef->size);
1497         break;
1498     case P_CHAR:
1499         if( !v->p_sdef )
1500             sprintf( text, "string: NULL" );
1501         else
1502             sprintf( text, "<%s>", v->p_char );
1503         break;
1504     case NUM:
1505         sprintf( text, "num: %d", v->num );
1506         break;
1507     case ASCII:
1508         sprintf( text, "ascii: '%s'", bin_to_ascii(v->ascii, 1) );
1509         break;
1510     }
1511   return buf;
1512 }
```

Listing 6.36. *main.c*—*main()* and a Symbol-Table-Printing Hook

```

1 #include <tools/debug.h>
2
3 int main( argc, argv )
4 char **argv;
5 {
6     UX( yy_init_occs(); ) /* Needed for yacc, called automatically by occs. */
7
8     argc = yy_get_args( argc, argv );
9
10    if( argc > 2 )           /* Generate trace if anything follows file name on */
11        enable_trace();      /* command line. */
12
13    yyparse();
14 }
15
16 /***** yyhook_a() *****/
17 yyhook_a()                  /* Print symbol table from debugger with ctrl-A. */
18 {                           /* Not used in yacc-generated parser. */
19     static int x = 0;
20     char buf[32];
21
22     sprintf( buf, "sym.%d", x++ );
23     yycomment( "Writing symbol table to %s\n", buf );
24     print_syms( buf );
25 }
26
27 yyhook_b()                  /* Enable/disable run-time trace with Ctrl-b. */
28 {                           /* Not used in yacc-generated parser. */
29     /* enable_trace() and disable_trace() are discussed */
30     /* below, when the gen() call is discussed. */
31     char buf[32];
32     if( yyprompt( "Enable or disable trace? (e/d): ", buf, 0 ) )
33     {
34         if( *buf == 'e' ) enable_trace();
35         else             disable_trace();
36     }
37 }
```

6.5 The Lexical Analyzer

Now that we've seen the tokens and the value-stack definition, we can look at the lexical analyzer specification (in Listing 6.37). This file is essentially the analyzer at the end of Appendix D, but it's been customized for the current parser by adding an attribute-passing mechanism. The **union** on lines 18 to 28, which reflects the **%union** definition in the *occs* input file, is used to pass attributes back to the parser. For example, when an **ASSIGNOP** is recognized on line 115, the first character of the lexeme is attached to the token by assigning a value to *yylval*, which is of the same type as a value-stack element. (The **union** definition on lines 18 to 28 doubles as an **extern** definition of *yylval*.) The parser pushes the contents of *yylval* onto the value stack when it shifts the state representing the token onto the state stack.

Also note that a newline is now recognized explicitly on line 132. The associated action prints a comment containing the input line number to the compiler's output file so that you can relate one to the other when you're debugging the output. You could also

use this action to pass input-line-number information to a debugger.

Listing 6.37. *c.lex*— C Lexical Analyzer

```

1  /* Lexical analyzer specification for C. This is a somewhat extended version
2   * of the one in Appendix D. The main difference is that it passes attributes
3   * for some tokens back to the parser, using the yylval mechanism to push
4   * the attribute onto the value stack.
5   */
6
7  %
8  #include <stdio.h>
9  #include <search.h>           /* Function prototype for bsearch(). */ */
10 #include <tools/debug.h>      /* Needed by symtab.h. */ */
11 #include <tools/hash.h>       /* Needed by symtab.h. */ */
12 #include <tools/l.h>
13 #include <tools/compiler.h>
14 #include "yyout.h"            /* Token defs. created by occs. Yacc uses y.tab.h. */ */
15 #include "symtab.h"           /* Needed to pass attributes to parser. */ */
16 #include "value.h"            /* ditto */ */
17
18 extern union {                /* This definition must duplicate the %union */
19     char      *p_char;        /* in c.y. */ */
20     symbol    *p_sym;
21     link      *p_link;
22     structdef *p_sdef;
23     specifier *p_spec;
24     value     *p_value;
25     int       integer;
26     int       ascii;
27 }
28 yyval;                      /* Declared by occs in yyout.c. */
29
30 extern FILE *yycodeout;      /* Declared by occs in yyout.c */
31
32 /*-----*/
33
34 #define YYERROR(text)    yyerror("%s\n", text); /* Does nothing in UNIX lex. */
35
36 /*-----*/
37 %
38
39 let    [_a-zA-Z]      /* Letter */ */
40 alnum [_a-zA-Z0-9]      /* Alphanumeric character */ */
41 h      [0-9a-fA-F]    /* Hexadecimal digit */ */
42 o      [0-7]          /* Octal digit */ */
43 d      [0-9]          /* Decimal digit */ */
44 suffix [UuLl]          /* Suffix in integral numeric constant */ */
45 white [\x00-\x09\x0b\s] /* White space: all control chars but \n */ */
46
47 %%
48 /* */
49 {
50     int i;
51
52     while( i = ii_input() )
53     {
54         if( i < 0 )
55             ii_flushbuf(); /* Discard lexeme. */

```

Listing 6.37. continued...

```

56                               else if( i == '*' && ii_lookahead(1) == '/' )
57                               {
58                                   ii_input();
59                                   break;                                /* Recognized comment.*/
60                               }
61
62
63                               if( i == 0 )
64                                   yyerror( "End of file in comment\n" );
65
66
67     "\"(\\".|[^\"])*\"      { return STRING; }
68
69     "\"(\\".|[^\"])*\n    yyerror("Adding missing \" to string constant\n");
70     yymore();
71
72     0{o}*{suffix}?          |
73     0x{h}+{suffix}?          |
74     [1-9]{d}*{suffix}?          return ICON ;
75
76     ({d}+|{d}+\.{d}*|{d}*\.{d}+)([eE][+-]?{d}+)?[fF]?   return FCON ;
77
78     "("                      return LP;
79     ")"                      return RP;
80     "{"                      return LC;
81     "}"                      return RC;
82     "["                      return LB;
83     "]"                      return RB;
84
85     "->"                   |
86     "."                      yyval.ascii = *yytext;
87                               return STRUCTOP;
88
89     "++"                     |
90     "--"                     yyval.ascii = *yytext;
91                               return INCOP;
92
93     [~!]                     yyval.ascii = *yytext;
94                               return UNOP;
95
96     "*"                      return STAR;
97
98     [/%]                     yyval.ascii = *yytext;
99                               return DIVOP;
100
101    "+"                      return PLUS;
102    "-"                      return MINUS;
103
104    <<|>>                  yyval.ascii = *yytext;
105                               return SHIFTOP;
106
107    [<>]=?                  yyval.ascii = yytext[1] ? (yytext[0]==>' ? 'G' : 'L')
108                                         : (yytext[0] );
109                               return RELOP;
110
111    [=!=]                    yyval.ascii = *yytext;
112                               return EQUOP;
113

```



Listing 6.37. continued...

Listing 6.37. continued...

```

174     { "void",      TYPE          },
175     { "while",     WHILE         }
176 };
177
178 static int      cmp( a, b )
179 KWORD   *a, *b;
180 {
181     return strcmp( a->name, b->name );
182 }
183
184 int      id_or_keyword( lex ) /* Do a binary search for a */
185 char    *lex;                /* possible keyword in Ktab */
186 {
187     KWORD      *p;           /* Return the token if it's */
188     KWORD      dummy;        /* in the table, NAME          */
189                           /* otherwise.                  */
190
191     dummy.name = lex;
192     p = bsearch( &dummy, Ktab, sizeof(Ktab)/sizeof(KWORD), sizeof(KWORD), cmp );
193
194     if( p )                  /* It's a keyword. */
195     {
196         yyval.ascii = *yytext;
197         return p->val;
198     }
199     else if( yyval.p_sym = (symbol *) findsym( Symbol_tab, yytext ) )
200         return (yyval.p_sym->type->tdef) ? TTYPE : NAME ;
201     else
202         return NAME;
203 }
```

6.6 Declarations

6.6.1 Simple Variable Declarations

This section moves from data structures to analyzing the code that manipulates those data structures. Declaration processing involves two main tasks: (1) you must assemble the linked lists that represent the types, attach them to symbols, and put the resulting structures into the symbol table, and (2) you must both generate C-code definitions for variables at fixed addresses and figure the frame-pointer-relative offsets for automatic variables. Many of the productions in the grammar work in concert to produce a single definition. Since these productions are discussed piecemeal, you should take a moment and look at the overall structure of the grammar in Appendix C. The declarations are all handled first, at the top of the grammar.

An example variable-declaration parse.

The best way to understand how the code generation works is to follow along as the parser works on an explicit example. (This is also a good approach for adding the code-generation actions in the first place: run a short sample through the parser observing the order in which reductions occur. Since actions at the far right of productions are executed in the same sequence as the reductions, you can see where the various actions need to be placed.) I'll use the following input to demonstrate simple declarations:

```
long int *x, y;
```

The parse of that input is shown, in symbolic form, in Table 6.13.¹⁴

Table 6.13. A Parse of `long int *x, y;`

	Stack	Next Action Taken by Parser
1	(empty)	Reduce by <code>ext_def_list → ε</code>
2	<code>ext_def_list</code>	Shift TYPE (long)
3	<code>ext_def_list TYPE</code>	Reduce by <code>typeSpecifier → TYPE</code>
4	<code>ext_def_list typeSpecifier</code>	Reduce by <code>typeOrClass → typeSpecifier</code>
5	<code>ext_def_list typeOrClass</code>	Reduce by <code>specifiers → typeOrClass</code>
6	<code>ext_def_list specifiers</code>	Shift TYPE (int)
7	<code>ext_def_list specifiers TYPE</code>	Reduce by <code>typeSpecifier → TYPE</code>
8	<code>ext_def_list specifiers typeSpecifier</code>	Reduce by <code>typeOrClass → typeSpecifier</code>
9	<code>ext_def_list specifiers typeOrClass</code>	Reduce by <code>specifiers → specifiers typeOrClass</code>
10	<code>ext_def_list specifiers</code>	Reduce by <code>optSpecifiers → specifiers</code>
11	<code>ext_def_list optSpecifiers</code>	Shift STAR
12	<code>ext_def_list optSpecifiers STAR</code>	Shift NAME (x)
13	<code>ext_def_list optSpecifiers STAR NAME</code>	Reduce by <code>newName → NAME</code>
14	<code>ext_def_list optSpecifiers STAR newName</code>	Reduce by <code>varDecl → newName</code>
15	<code>ext_def_list optSpecifiers STAR varDecl</code>	Reduce by <code>varDecl → STAR varDecl</code>
16	<code>ext_def_list optSpecifiers varDecl</code>	Reduce by <code>extDecl → varDecl</code>
17	<code>ext_def_list optSpecifiers extDecl</code>	Reduce by <code>extDeclList → extDecl</code>
18	<code>ext_def_list optSpecifiers extDeclList</code>	Shift COMMA
19	<code>ext_def_list optSpecifiers extDeclList COMMA</code>	Shift NAME (y)
20	<code>ext_def_list optSpecifiers extDeclList COMMA NAME</code>	Reduce by <code>newName → NAME</code>
21	<code>ext_def_list optSpecifiers extDeclList COMMA newName</code>	Reduce by <code>varDecl → newName</code>
22	<code>ext_def_list optSpecifiers extDeclList COMMA varDecl</code>	Reduce by <code>extDecl → varDecl</code>
23	<code>ext_def_list optSpecifiers extDeclList COMMA extDecl</code>	Reduce by <code>extDeclList → extDeclList COMMA extDecl</code>
24	<code>ext_def_list optSpecifiers extDeclList</code>	Reduce by <code>{3} → ε</code>
25	<code>ext_def_list optSpecifiers extDeclList {3}</code>	Shift SEMI
26	<code>ext_def_list optSpecifiers extDeclList {3} SEMI</code>	Reduce by <code>extDef → optSpecifiers extDeclList {3} SEMI</code>
27	<code>ext_def_list extDef</code>	Reduce by <code>extDefList → extDefList extDef</code>
28	<code>ext_def_list</code>	Reduce by <code>program → extDefList (accept)</code>

Since all actions are performed as part of a reduction, the best way to see how the code works is to follow the order of reductions during the parse. First, notice how the ϵ production in

```
extDefList
: extDefList extDef
| /* epsilon */
```

(on line 189 of Listing 6.38) is done first. This is always the case in left-recursive list productions: the nonrecursive component (whether or not it's ϵ) must be reduced first in order to put the recursive left-hand side onto the stack.¹⁵ In the current example, the

14. If you have the distribution disk mentioned in the Preface, you can use the visible parser to see the parse process in action. The file `c.exe` is an executable version of the compiler described in this chapter. Get the parse started by creating a file called `test.c` containing the single line:

```
long int *x, y;
```

Then invoke the parser with the command: `c test.c`. Use the space bar to singlestep through the parse.

15. The nonrecursive element of a right-recursive list production is always done last.

parser must get an *ext_def_list* onto the stack in order to be able to reduce by $\text{ext_def_list} \rightarrow \text{ext_def_list ext_def}$ at some later time, and the only way to get that initial *ext_def_list* onto the stack is to reduce by $\text{ext_def_list} \rightarrow \epsilon$.

Listing 6.38. c.y—Initialization and Cleanup Productions

```

183  %%
184  program : ext_def_list { clean_up(); }
185  ;
186
187 ext_def_list
188   : ext_def_list ext_def
189   | /* epsilon */
190   {
191     yydata( "#include <tools/virtual.h>\n" );
192     yydata( "#define T(x)\n" );
193     yydata( "SEG(data)\n" );
194     yycode( "\nSEG(code)\n" );
195     yybss ( "\nSEG(bss)\n" );
196   }
197 ;

```

Initializations done in
nonrecursive list element.

Generate SEG directives.

Specifier processing.

Syntactically, a program is a list of external definitions because of the productions on lines 187 to 189 of Listing 6.38. A reduction by the ϵ production on line 189 is always the first action taken by the parser, regardless of the input. This behavior is exploited in the current compiler to do various initializations on lines 189 to 196 of Listing 6.38. The appropriate SEG directives are generated at the tops of the various segments, and a **#include** for `<tools/virtual.h>` is output at the top of the data segment, which is at the top of the output file after the three segments are combined at the end of the compilation. I'll discuss the empty `T()` macro on line 192 in a moment.

The associated cleanup actions are all done in the previous production:

```
program : ext_def_list { clean_up(); }
```

on line 184. Remember, this is a bottom-up parser, so the reduction to the goal symbol is the last action taken by the parser. The `clean_up()` action coalesces the output streams and deletes temporary files.

After initializing via the ϵ production, the parser starts processing the specifier component of the declaration. All of the specifier productions are shown together in Listing 6.39. Three types of specifier lists are supported:

<code>opt_specifiers</code>	Zero or more types and storage classes mixed together.
<code>specifiers</code>	One or more types and storage classes mixed together.
<code>type</code>	One or more types. (No storage classes are permitted.)

Note that the parser is not checking semantics here. It just recognizes collections of possible types and storage classes without testing for illegal combinations like a **short long**.

The parser starts by shifting the **TYPE** token (**long** in the current input) and then reduces $\text{type_specifier} \rightarrow \text{TYPE}$ (on line 229 of Listing 6.39). The associated action calls `new_type_spec()`, which gets a new link and initializes it to a specifier of the correct type. A pointer to the link structure is attached to the `type_specifier` by assigning it to `$$`. From here on out, that pointer is on the value stack at the position corresponding to the `type_specifier`. If a storage class is encountered instead of a **TYPE**, the action on line 226 of Listing 6.39 is executed instead of the action on line 189. This action modifies the storage-class component of the link rather than the type, but is

Pointer to `link` used as
an attribute.

Listing 6.39. *c.y*—Specifiers

```

198 opt_specifiers
199      : CLASS TTYPE { set_class_bit( 0, $2->etype ); /* Reset class. */ /
200      | set_class_bit( $1, $2->etype ); /* Add new class. */
201      | $2 = $2->type ;
202      }
203      | TTYPE { set_class_bit( 0, $1->etype ); /* Reset class bits. */
204      | $2 = $1->type ;
205      }
206      | specifiers
207      | /* empty */
208      {
209      |           %%prec COMMA
210      |           $2 = new_link();
211      |           $2->class = SPECIFIER;
212      |           $2->NOUN = INT;
213      }
214 specifiers
215      : type_or_class
216      | specifiers type_or_class { spec_cpy( $$, $2 );
217      |                         discard_link_chain( $2 ); }
218      ;
219 type
220      : typeSpecifier
221      | type typeSpecifier { spec_cpy( $$, $2 );
222      |                         discard_link_chain( $2 ); }
223      ;
224 type_or_class
225      : typeSpecifier
226      | CLASS { $$ = new_class_spec( $1 ); }
227      ;
228 typeSpecifier
229      : TYPE { $$ = new_type_spec( yytext ); }
230      | enumSpecifier { $$ = new_type_spec( "int" ); }
231      | structSpecifier { $$ = new_link();
232      |                     $2->class = SPECIFIER;
233      |                     $2->NOUN = STRUCTURE;
234      |                     $2->V_STRUCTURE = $1;
235      }
236      ;

```

otherwise the same as the earlier one. Both `new_type_spec()` and `new_class_spec()` are in Listing 6.40.

The next reductions are $\text{type_or_class} \rightarrow \text{type_specifier}$ and $\text{specifiers} \rightarrow \text{type_or_class}$, neither of which has an associated action. The pointer-to-link attribute is carried along with each reduction because of the default $$$=\1 action that is supplied by the parser. The pointer to the link created in the initial reduction is still on the value stack, but now at the position corresponding to the *specifiers* nonterminal.

The parser now processes the `int`. It performs the same set of reductions that we just looked at, but this time it reduces by

$\text{specifiers} \rightarrow \text{specifiers type_or_class}$

rather than by

$\text{specifiers} \rightarrow \text{type_or_class}$

as was done earlier. (We're on line nine of Table 6.13. The associated action is on lines

$\text{type_or_class} \rightarrow \text{type_specifier}$
 $\text{specifiers} \rightarrow \text{type_or_class}$

Listing 6.40. *decl.c*—Create and Initialize a link

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/debug.h>
4 #include <tools/hash.h>
5 #include <tools/l.h>
6 #include <tools/compiler.h>
7 #include <tools/c-code.h>
8
9 #include "symtab.h"
10 #include "value.h"
11 #include "proto.h"
12
13 /* DECL.C      This file contains support subroutines for those actions in c.y
14 *           that deal with declarations.
15 */
16
17 extern void yybss(), yydata();
18
19 /*-----*/
20
21 PUBLIC link *new_class_spec( first_char_of_lexeme )
22 int first_char_of_lexeme;
23 {
24     /* Return a new specifier link with the sclass field initialized to hold
25      * a storage class, the first character of which is passed in as an argument
26      * ('e' for extern, 's' for static, and so forth).
27     */
28
29     link *p = new_link();
30     p->class = SPECIFIER;
31     set_class_bit( first_char_of_lexeme, p );
32     return p;
33 }
34
35 /*-----*/
36
37 PUBLIC void set_class_bit( first_char_of_lexeme, p )
38 int first_char_of_lexeme;
39 link *p;
40 {
41     /* Change the class of the specifier pointed to by p as indicated by the
42      * first character in the lexeme. If it's 0, then the defaults are
43      * restored (fixed, nonstatic, nonexternal). Note that the TYPEDEF
44      * class is used here only to remember that the input storage class
45      * was a typedef, the tdef field in the link is set true (and the storage
46      * class is cleared) before the entry is added to the symbol table.
47     */
48
49     switch( first_char_of_lexeme )
50     {
51     case 0:   p->SCLASS = FIXED;
52             p->STATIC = 0;
53             p->EXTERN = 0;
54             break;
55
56     case 't': p->SCLASS = TYPEDEF ; break;
57     case 'r': p->SCLASS = REGISTER ; break;

```



Listing 6.40. continued...

```

58     case 's': p->STATIC = 1      ; break;
59     case 'e': p->EXTERN = 1      ; break;
60
61     default : yyerror("INTERNAL, set_class_bit: bad storage class '%c'\n",
62                         first_char_of_lexeme);
63                         exit( 1 );
64     break;
65 }
66 }
67
68 /*-----*/
69
70 PUBLIC link *new_type_spec( lexeme )
71 char *lexeme;
72 {
73     /* Create a new specifier and initialize the type according to the indicated
74      * lexeme. Input lexemes are: char const double float int long short
75      *                                signed unsigned void volatile
76      */
77
78     link *p = new_link();
79     p->class = SPECIFIER;
80
81     switch( lexeme[0] )
82     {
83     case 'c': if( lexeme[1]=='h' )           /* char | const      */
84                     p->NOUN = CHAR ;          /* (Ignore const.) */
85                     break;
86     case 'd':                           /* double           */
87     case 'f': yyerror("No floating point\n"); /* float            */
88                     break;
89
90     case 'i': p->NOUN      = INT;    break;   /* int              */
91     case 'l': p->LONG       = 1;      break;   /* long             */
92     case 'u': p->UNSIGNED   = 1;      break;   /* unsigned         */
93
94     case 'v': if( lexeme[2] == 'i' )        /* void | volatile */
95                     p->NOUN = VOID;          /* ignore volatile */
96                     break;
97     case 's': break;                      /* short | signed */
98     }
99
100    return p;
101 }

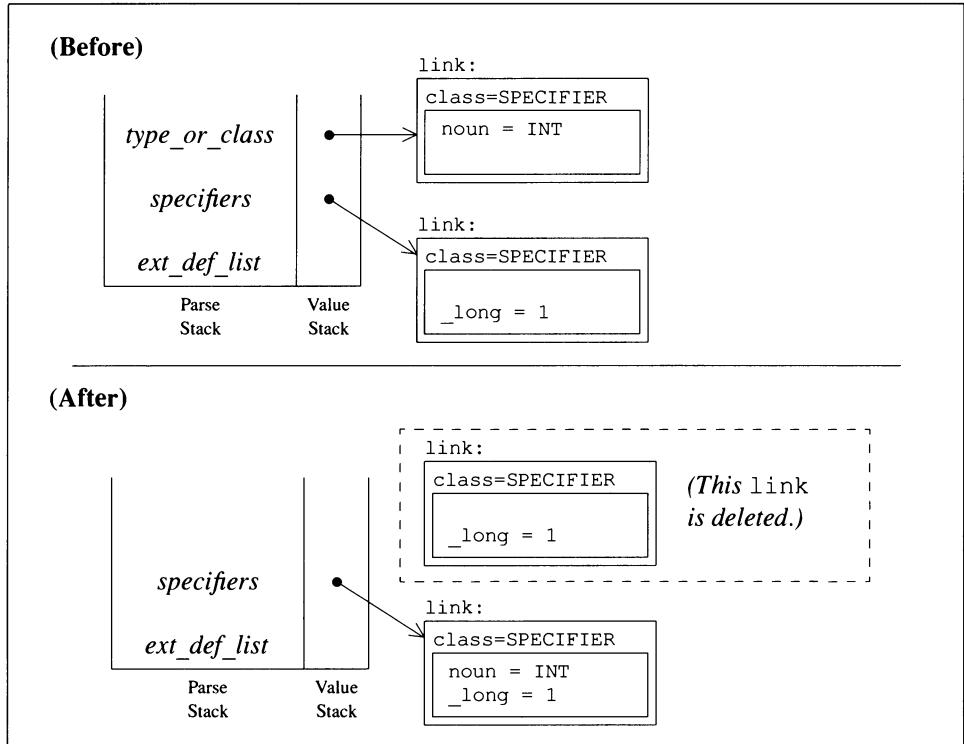
```

216 to 217 of Listing 6.39.) There are currently three attributes of interest: \$1 and \$\$ (which is initialized to \$1 by the parser before executing the action) hold a pointer to the link structure that was assembled when the **long** was processed; \$2 points at the link for the new list element—the **int**. The action code merges the two structures by copying all relevant fields from \$2 to \$\$\$. The second link (\$2) is then discarded. The action is illustrated in Figure 6.14.

If additional types or storage classes are present in the input, the parser loops some more, creating new links for each new list element and merging those new links into the existing one, discarding the new link after it's merged.

Having collected the entire specifier list, the parser now reduces by *opt_specifiers* *opt_specifiers→specifiers* on line 206 of Listing 6.40. There is no action. If the

Figure 6.14. Merging Links



TTYPE, processing a **typedef**.

declaration were in terms of a **typedef**, *opt_specifiers* → **TTYPE** (just above the current production on line 203 of Listing 6.40) is executed instead of the list-collecting reductions we just looked at. The scanner attaches an attribute to the **TTYPE** token: a pointer to a symbol structure representing the **typedef**. The name field in the symbol is the type name rather than a variable name. The action on line 204 of Listing 6.40 passes the type string for this symbol back up the stack as *opt_specifiers*' attribute. This type chain differs from a normal one only in that the tdef bit is set in the leftmost link in the chain. The parser needs this information because a type that comes from a **typedef** could be an entire chain, including the declarators. A non-**typedef** declaration always yields a single specifier. If no specifier at all is present, then the action on lines 207 to 212 of Listing 6.40 is executed.¹⁶ This action just sets up things as if an **int** had been found in the input.

The parser shelves the specifier on the value stack for a moment to move on to the declarator component of the declaration—it stores a pointer to the accumulated specifier link on the value stack as an attribute attached to the most recently pushed *specifier* nonterminal. Declarator processing starts with a shift of the **STAR** and **NAME** tokens on lines 11 and 12 of Table 6.13 on page 523. The parser then starts reducing by the productions in Listing 6.41, below.

The first reduction is *new_name* → **NAME**. The action is on line 282 of Listing 6.41. There are two name productions to take care of two different situations in which names

Declarator processing.

Identifier processing in declaration,
new_name → **NAME**.

¹⁶ *x;*, with no explicit type, is a perfectly legitimate declaration, the default **int** type and **extern** storage classes are used.

Listing 6.41. c.y— Variable Declarators

```

237 var_decl
238     : new_name           %prec COMMA /* This production is done first. */
239
240     | var_decl LP RP      { add_declarator( $$, FUNCTION ); }
241     | var_decl LP var_list RP { add_declarator( $$, FUNCTION );
242                               discard_symbol_chain( $3 );
243                               }
244     | var_decl LB RB
245     {
246         /* At the global level, this must be treated as an array of
247          * indeterminate size; at the local level this is equivalent to
248          * a pointer. The latter case is patched after the declaration
249          * is assembled.
250         */
251
252         add_declarator( $$, ARRAY );
253         $$->etype->NUM_ELE = 0;
254
255         YYD( yycomment("Add POINTER specifier\n") );
256     }
257
258     | var_decl LB const_expr RB
259     {
260         add_declarator( $$, ARRAY );
261         $$->etype->NUM_ELE = $3;
262
263         YYD(yycomment("Add array[%d] spec.\n", $$->etype->NUM_ELE));
264     }
265     | STAR var_decl          %prec UNOP
266     {
267         add_declarator( $$ = $2, POINTER );
268         YYD( yycomment("Add POINTER specifier\n") );
269     }
270
271     | LP var_decl RP { $$ = $2; }
272     ;
273
274 /*
275  * Name productions. new_name always creates a new symbol, initialized with the
276  * current lexeme. Name returns a preexisting symbol with the associated name
277  * (if there is one); otherwise, the symbol is allocated. The NAME token itself
278  * has a NULL attribute if the symbol doesn't exist, otherwise it returns a
279  * pointer to a "symbol" for the name.
280 */
281
282 new_name: NAME { $$ = new_symbol( yytext, Nest_lev ); }
283 ;
284
285 name   : NAME { if( !$1 || $1->level != Nest_lev )
286                  $$ = new_symbol( yytext, Nest_lev );
287                }
288 ;

```

are used. Normally, the scanner looks up all identifiers in the symbol table to see whether or not they represent types so that it can return a **TTYPE** token when appropriate. In order to avoid unnecessary lookups, the scanner attaches a symbol-table pointer for an existing identifier to the **NAME** token if the identifier is in the table, otherwise the

Interaction between
scanner and code gen-
erator via symbol table.

name → **NAME**.

Nest_lev.

attribute is NULL. This way the table lookup doesn't have to be done a second time in the code-generation action. The *new_name* takes care of newly-created identifiers. It ignores this passed-back symbol pointer, and allocates a new symbol structure, passing a pointer to the symbol back up the parse tree as an attribute attached to the *new_name* nonterminal. The *name* → **NAME** action (on lines 285 to 286, of Listing 6.41) uses the existing symbol. The test on line 285 identifies whether the returned symbol is at the current scoping level. If it is not, then the parser assumes that this is a new symbol (such as a local variable) that just happens to have the same name as an existing one, and allocates a new symbol structure. *Nest_lev* keeps track of the current scope level; it is declared in Listing 6.42, which is part of the occs-input-file definitions section.

Listing 6.42. *c.y*— Identify Nesting Level (from Occs Definitions Section)

```
119 %{
120 int     Nest_lev;      /* Current block-nesting level.          */
121 %}
```

var_dcl → *new_name*.

The next reduction, on line 15 of Table 6.13 on page 523), is *var_dcl* → *new_name*. The production is on line 238 of Listing 6.42. The only action is the implicit $\$\$=\$1$, which causes the symbol pointer returned from *new_name* to be passed further up the tree, attached to the *var_dcl*. The *%prec* at the right of the line eliminates a shift/reduce conflict by assigning a very low precedence to the current production—the technique is discussed in Appendix E.

var_decl → **STAR**
 var_decl.
add_declarator().

The parser reduces by *var_decl* → **STAR** *var_decl* next. (We're moving from line 15 to line 16 of Table 6.13 on page 523—the action is on line 265 of Listing 6.42.) The action calls *add_declarator()* to add a pointer-declarator link to the type chain in the symbol that was created when the name was processed. The process is similar to that used to assemble an NFA in Chapter Two. Figure 6.15 shows the parse and value stacks both before and after this reduction is performed.

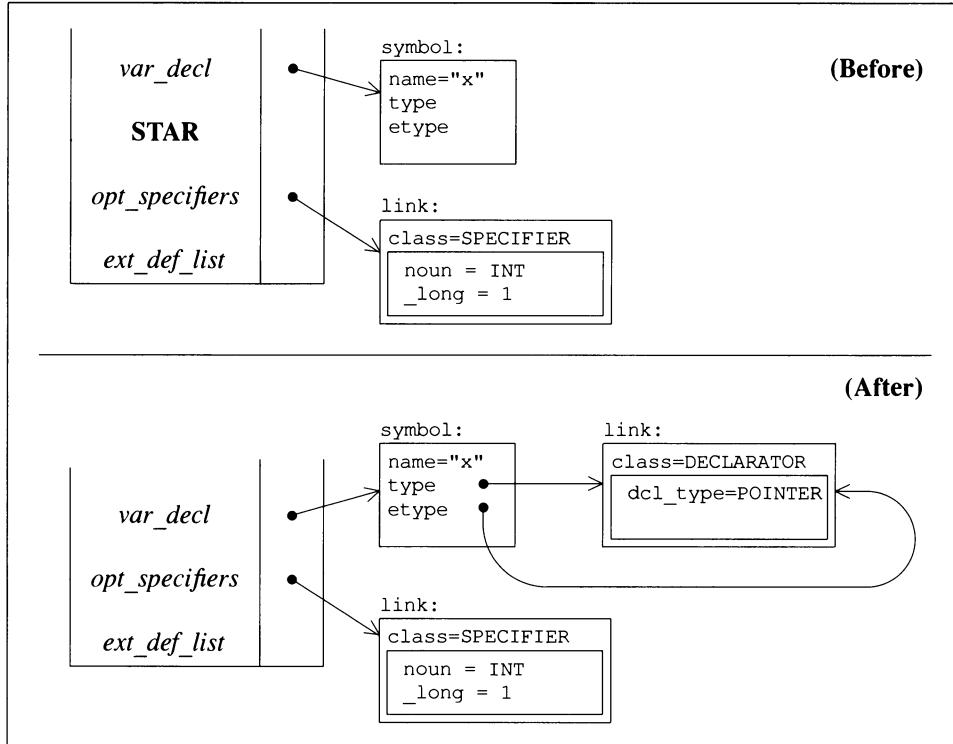
Note that the synthesized attribute for every right-hand side of this production is the same symbol structure that was allocated when the name was processed. The actions just add link structures to the symbol's type chain.

var_decl

If the declaration had been **int** **x*, both stars would have been shifted initially, and the current reduction would have executed twice in succession, thereby adding a second pointer-declarator link to the end of the type chain in the symbol structure. The other productions that share the *var_decl* left-hand side do similar things, adding declarator links either for arrays or for functions, as appropriate, to the end of the type chain in the current symbol structure. Note that the actions on lines 240 to 243 handle the “function” component of a function-pointer declaration. I'll look at function declarations, including the *var_list* nonterminal, in a moment—briefly, the *var_list* takes care of function prototypes; the associated attribute is a pointer to the head of a linked list of symbol structures, one for each function argument. The action here just discards all the symbols in the list. Similarly, *const_expr* (on line 258) handles integer-constant expressions. This production is also discussed below, but the associated attribute is an integer. I'm using the <num> field in the value-stack union to hold the value.

Dummy *const_expr* production used for grammar development.

As an aside, a minor problem came up when adding the actions to the grammar—you can't do the expression-processing actions until the declarations are finished, but you need to use a constant expression to process an array declaration. I solved the problem by providing a dummy action of the form:

Figure 6.15. Adding Declarators to the Type Chain

```
const_expr : expr { $$ = 10; }
;
```

until the declarations were working. The action was later replaced with something more reasonable.

When the parser finishes with the declarator elements (when the lookahead is a **COMMA**), the type chain in the `symbol` structure holds a linked list of declarator links, and the specifier is still on the value stack at a position corresponding to the `opt_specifiers` nonterminal. The comma is shifted, and the parser goes through the entire declarator-processing procedure again for the `y` (on lines 20 to 23 of Table 6.13 on page 523).

Now the parser starts to create the cross links for symbol-table entry—the links that join declarations for all variables at the current scoping level. It does this using the productions in Listing 6.43. The first reduction of interest is $\text{ext_decl_list} \rightarrow \text{ext_decl}$ executed on line 17 of Table 6.13. The associated action, on line 298 of Listing 6.43, puts a NULL pointer onto the value stack. This pointer marks the end of the linked list. The parse proceeds as just described, until the declaration for `y` has been processed, whereupon the parser links the two declarations together. The parse and value stacks, just before and just after the reduction by:

$\text{ext_decl_list} \rightarrow \text{ext_decl_list} \text{ COMMA ext_decl}$

are shown in Figure 6.16. and the code that does the linking is on lines 308 to 313 of Listing 6.43. If there were more comma-separated declarators in the input, the process would continue in this manner, each successive element being linked to the head of the list in turn.

Create cross links.
 $\text{ext_decl_list} \rightarrow \text{ext_decl_list}$.

$\text{ext_decl_list} \rightarrow$
 ext_decl_list
 $\text{COMMA ext_decl}.$

Listing 6.43. *c.y*—Function Declarators

```

289 /*-----  
290 * Global declarations: take care of the declarator part of the declaration.  
291 * (The specifiers are handled by specifiers).  
292 * Assemble the declarators into a chain, using the cross links.  
293 */  
294  
295 ext_decl_list  
296     : ext_decl  
297     {  
298         $$->next = NULL;                                /* First link in chain. */  
299     }  
300     | ext_decl_list COMMA ext_decl  
301     {  
302         /* Initially, $1 and $$ point at the head of the chain.  
303         * $3 is a pointer to the new declarator.  
304         */  
305  
306         $3->next = $1;  
307         $$      = $3;  
308     }  
309 ;  
310  
311 ext_decl  
312     : var_decl  
313     | var_decl EQUAL initializer { $$->args = (symbol *)$3; }  
314     | funct_decl  
315 ;  
316

```

Initializers, *symbol.args*.

The only other issue of interest is the initializer, used on line 313 of Listing 6.43. I'll defer discussing the details of initializer processing until expressions are discussed, but the attribute associated with the initializer is a pointer to the head of a linked list of structures that represent the initial values. The *args* field of the *symbol* structure is used here to remember this pointer. You must use a cast to get the types to match.

Function declarators,
funct_decl.

Before proceeding with the sample parse, it's useful to back up a notch and finish looking at the various declarator productions. There are two types of declarators not used in the current example, function declarators and abstract declarators. The function-processing productions start in Listing 6.44. First, notice the *funct_decl* productions are almost identical to the *var_decl* productions that were examined earlier. They both assemble linked lists of declarator links in a *symbol* structure that is passed around as an attribute. The only significant additions are the right-hand sides on lines 329 to 341, which handle function arguments.

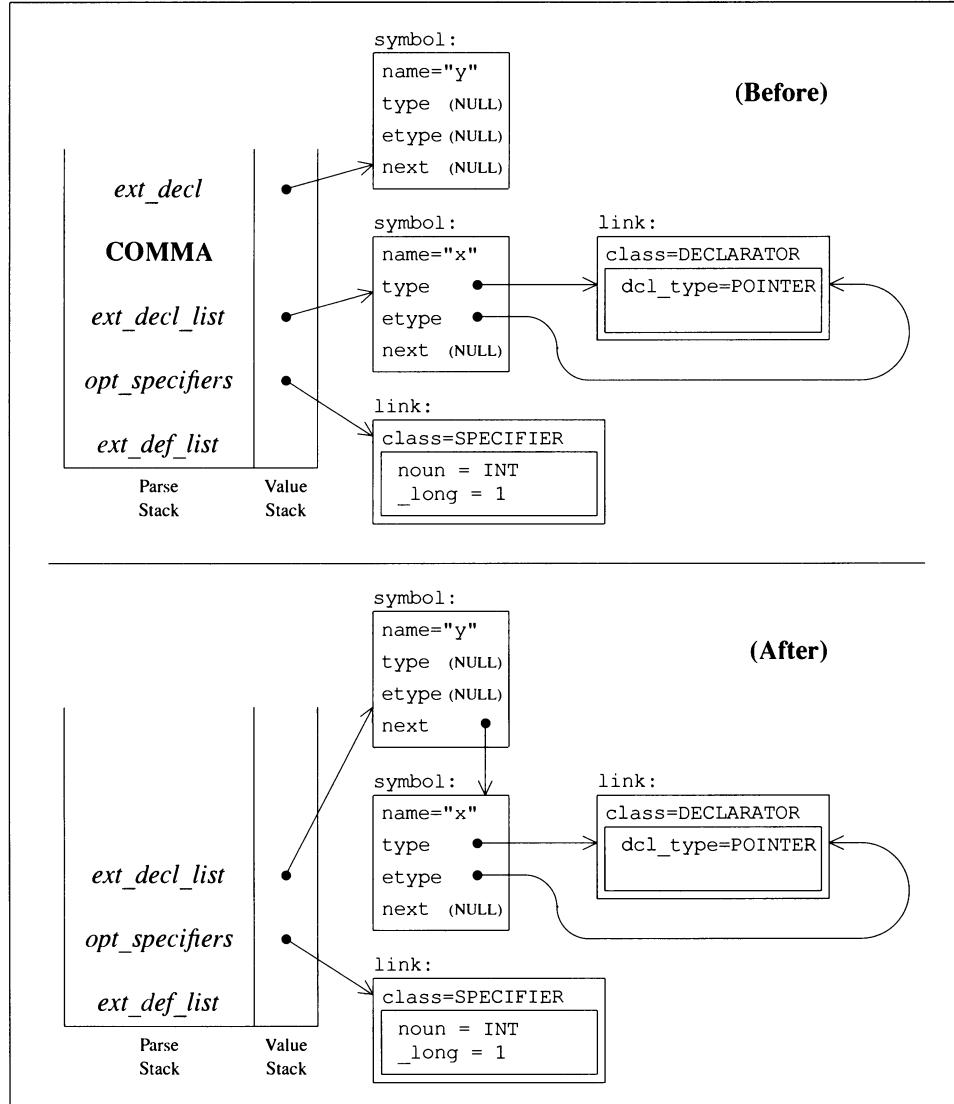
Function-argument de-
clarations.

The same *funct_decl* productions are used both for function declarations (**externs** and prototypes) and function definitions (where a function body is present)—remember, these productions are handling only the declarator component of the declaration. The situation is simplified because prototypes are ignored. If they were supported, you'd have to detect semantic errors such as the following one in which the arguments don't have names. The parser accepts the following input without errors:

```

foo( int, long )
{
    /* body */
}

```

Figure 6.16. A Reduction by $\text{ext_decl_list} \rightarrow \text{ext_decl_list} \text{ COMMA } \text{ext_decl}$ 

The `Nest_lev` variable is modified in the imbedded actions on lines 329 and 335 because function arguments are actually at the inner scoping level.

The function arguments can take two forms. A `name_list` is a simple list of names, as `name_list`, used for arguments in the older, K&R style, declaration syntax:

```
hobbit( frito, bilbo, spam )
short frito, bilbo;
```

A `var_list` takes care of both the new, C++-style syntax:

```
hobbit( short frito, short bilbo, int spam );
```

and function prototypes that use abstract declarators (declarations without names) such as the following:

```
hobbit( short, short, int );
```

All of these forms are recognized by the parser, but the last one is just ignored in the

Listing 6.44. c.y—Function Declarators

```

317 funct_decl
318     : STAR funct_decl          { add_declarator( $$ = $2 , POINTER ); }
319     | funct_decl LB RB        { add_declarator( $$, ARRAY );
320                               $$->etype->NUM_ELE = 0;
321                           }
322     | funct_decl LB const_expr RB { add_declarator( $$, ARRAY );
323                                   $$->etype->NUM_ELE = $3;
324                           }
325     | LP funct_decl RP         { $$ = $2;
326     | funct_decl LP RP        { add_declarator( $$, FUNCTION );
327     | new_name LP RP           { add_declarator( $$, FUNCTION );
328                           }
329     | new_name LP { ++Nest_lev; } name_list { --Nest_lev; } RP
330     {
331         add_declarator( $$, FUNCTION );
332
333         $4           = reverse_links( $4 );
334         $$->args = $4;
335     }
336     | new_name LP { ++Nest_lev; } var_list { --Nest_lev; } RP
337     {
338         add_declarator( $$, FUNCTION );
339         $$->args = $4;
340     }
341 ;
342 name_list
343     : new_name                 {
344         $$->next      = NULL;
345         $$->type       = new_link();
346         $$->type->class = SPECIFIER;
347         $$->type->SCLASS = AUTO;
348     }
349     | name_list COMMA new_name {
350         $$           = $3;
351         $$->next      = $1;
352         $$->type       = new_link();
353         $$->type->class = SPECIFIER;
354         $$->type->SCLASS = AUTO;
355     }
356 ;
357 var_list
358     : param_declaration        { if($1) $$->next = NULL; }
359     | var_list COMMA param_declaration { if($3)
360                                         {
361                                             $$           = $3;
362                                             $3->next = $1;
363                                         }
364                                         }
365 ;
366 param_declaration
367     : type var_decl            { add_spec_to_decl($1, $$ = $2 ); }
368     | abstract_decl             { discard_symbol ($1); $$ = NULL; }
369     | ELLIPSIS                  { $$ = NULL; }
370 ;

```

present application: The *abstract_decl* production used on line 368 of Listing 6.44 creates a symbol structure, just like the ones we've been discussing. There's no name, however. The action on line 368 just throws it away. Similarly, the **ELLIPSIS** that's used for ANSI variable-argument lists is just ignored on the next line.

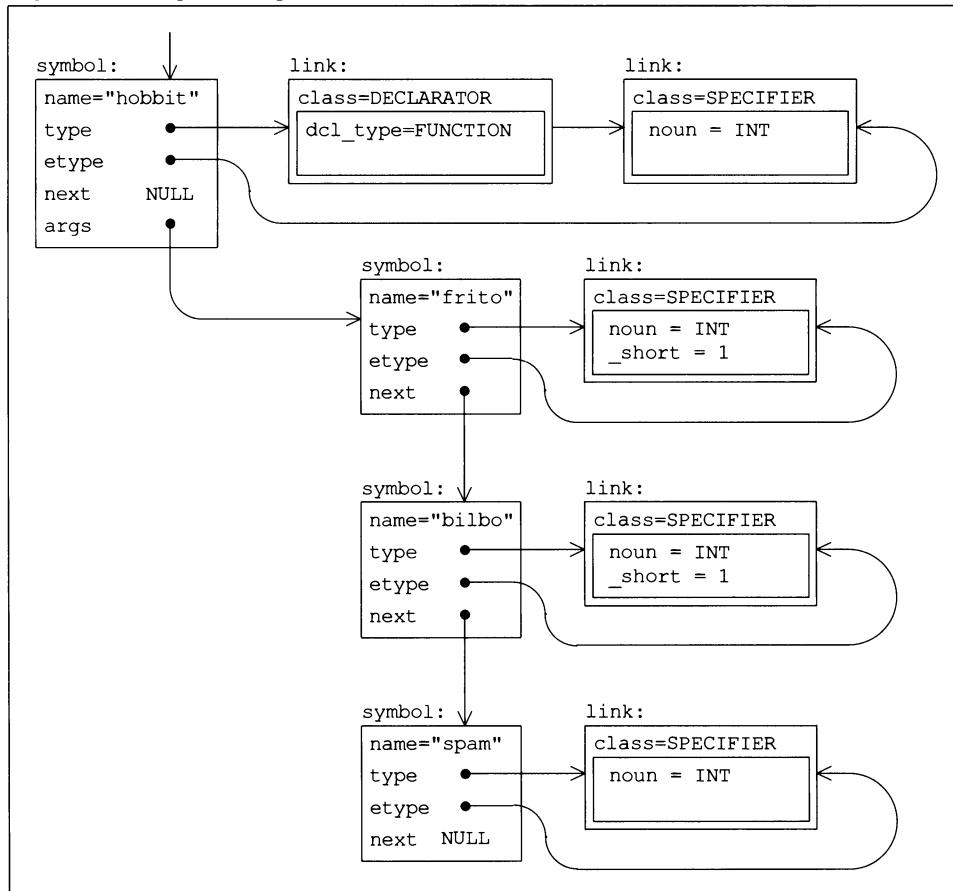
Abstract declarators,
abstract_decl.

The *name_list* and *var_list* productions are found on lines 342 to 365 of Listing 6.44. They differ from one another in only one significant way: A *name_list* creates a linked list of symbols, one for each name in the list, and the type chains for all these symbols are single specifier links representing **ints**. A *var_list* takes the types from the declarations rather than supplying an **int** type. By the time the parser gets back up to lines 329 or 336 of Listing 6.45-1 the argument list will have been processed, and the attribute attached to the *name_list* or *var_list* will be a pointer to the head of a linked list of symbol structures. The `$$->args = $4` attaches this list to the `args` field of the symbol that represents the function itself. Figure 6.17 shows the way that

```
hobbit( short frito, short bilbo, int spam );
```

is represented once all the processing is finished.

Figure 6.17. Representing `hobbit(short frito, short bilbo, int spam)`



Note that the list of arguments is assembled in reverse order because, though the list is processed from left to right, each new element is added to the head of the list. The `reverse_links()` call on line 333 of Listing 6.44 goes through the linked list of symbols and reverses the direction of the `next` pointers. It returns a pointer to the new head of the chain (formerly the end of the chain).

Argument declarations
assembled in reverse
order,
`reverse_links()`.

Merge declarator and specifier,
`add_spec_to_decl()`.

Abstract declarators.

One other subroutine is of particular interest: `add_spec_to_decl()`, called on line 367 of Listing 6.44, merges together the specifier and declarator components of a declaration. It's shown in Listing 6.45, below. Passed a pointer to a link that represents the specifier, and a pointer `symbol` that contains a type chain representing the declarator, it makes a copy of the specifier link [with the `clone_type()` call on line 132] and tacks the copy onto the end of the type chain.

The last kind of declarator in the grammar is an abstract declarator, handled in Listing 6.46, below. If abstract declarators were used only for declarations, the productions on lines 371 to 388 could be devoid of actions. You need abstract declarators for the cast operator and `sizeof` statement, however. The actions here work just like all the other declarator productions; the only difference is that the resulting `symbol` attribute has a type but no name. The `symbol` structure is allocated on line 380—the ϵ production takes the place of the identifier in the earlier declarator productions.

Returning to the parse in Table 6.13 on page 523, the parser has just finished with the `ext_decl_list` and is about to reduce by $\{3\} \rightarrow \epsilon$. This production is supplied by `occs` to process the imbedded action on lines 390 to 401 of Listing 6.47, below. `Occts` translates:

```
ext_def : opt_specifiers ext_decl_list {action...} SEMI
;
```

as follows:

```
ext_def : opt_specifiers ext_decl_list {3} SEMI
;
{3}    : /* empty */ {action...}
;
```

so that it can do the imbedded action as part of a reduction.

The action does three things: it merges the specifier and declarator components of the declaration, puts the new declarations into the symbol table, and generates the actual declarations in the output. The action must precede the `SEMI` because of a problem caused by the way that the parser and lexical analyzer interact with one another. The lexical analyzer uses the symbol table to distinguish identifiers from the synthetic types created by a `typedef`, but symbol-table entries are also put into the symbol table by the current production. The problem is that the input token is used as a lookahead symbol. It can be read well in advance of the time when it is shifted, and several reductions can occur between the read and the subsequent shift. In fact, the lookahead is required to know which reductions to perform. So the next token is always read immediately after shifting the previous token. Consider the following code:

```
typedef int itype;
itype      x;
```

If the action that adds the new type to the symbol table followed the `SEMI` in the grammar, the following sequence of actions would occur:

- Shift the `SEMI`, and read the lookahead symbol. `itype` has not been added to the symbol table yet, so the scanner returns a `NAME` token.
- Reduce by $ext_def \rightarrow \{opt_specifiers\ ext_decl_list\} \ SEMI$, adding the `itype` to the symbol table.

The problem is solved by moving the action forward in the production, so that the parser correctly acts as follows:

Problems with scanner,
code-generator interaction

Listing 6.45. *decl.c*— Add a Specifier to a Declaration

```
102 void add_spec_to_decl( p_spec, decl_chain )
103 link *p_spec;
104 symbol *decl_chain;
105 {
106     /* p_spec is a pointer either to a specifier/declarator chain created
107      * by a previous typedef or to a single specifier. It is cloned and then
108      * tacked onto the end of every declaration chain in the list pointed to by
109      * decl_chain. Note that the memory used for a single specifier, as compared
110      * to a typedef, may be freed after making this call because a COPY is put
111      * into the symbol's type chain.
112      *
113      * In theory, you could save space by modifying all declarators to point
114      * at a single specifier. This makes deletions much more difficult, because
115      * you can no longer just free every node in the chain as it's used. The
116      * problem is complicated further by typedefs, which may be declared at an
117      * outer level, but can't be deleted when an inner-level symbol is
118      * discarded. It's easiest to just make a copy.
119      *
120      * typedefs are handled like this: If the incoming storage class is TYPEDEF,
121      * then the typedef appeared in the current declaration and the tdef bit is
122      * set at the head of the cloned type chain and the storage class in the
123      * clone is cleared; otherwise, the clone's tdef bit is cleared (it's just
124      * not copied by clone_type().
125      */
126
127 link *clone_start, *clone_end ;
128 link **p;
129
130 for( ; decl_chain ; decl_chain = decl_chain->next )
131 {
132     if( !(clone_start = clone_type(p_spec, &clone_end)) )
133     {
134         yyerror("INTERNAL, add_typedef_: Malformed chain (no specifier)\n");
135         exit( 1 );
136     }
137     else
138     {
139         if( !decl_chain->type )                                /* No declarators. */
140             decl_chain->type = clone_start ;
141         else
142             decl_chain->etype->next = clone_start;
143
144             decl_chain->etype = clone_end;
145
146         if( IS_TYPEDEF(clone_end) )
147         {
148             set_class_bit( 0, clone_end );
149             decl_chain->type->tdef = 1;
150         }
151     }
152 }
153 }
```

Listing 6.46. *c.y*— Abstract Declarators

```

371 abstract_decl
372     : type abs_decl { add_spec_to_decl ( $1, $$ = $2 ); }
373     | TTYPE abs_decl {
374         $$ = $2;
375         add_spec_to_decl( $1->type, $2 );
376     }
377 ;
378
379 abs_decl
380     : /* epsilon */ { $$ = new_symbol("",0); }
381     | LP abs_decl RP LP RP { add_declarator( $$ = $2, FUNCTION ); }
382     | STAR abs_decl { add_declarator( $$ = $2, POINTER ); }
383     | abs_decl LB RB { add_declarator( $$, POINTER ); }
384     | abs_decl LB const_expr RB { add_declarator( $$, ARRAY );
385         $$->etype->NUM_ELE = $3;
386     }
387     | LP abs_decl RP { $$ = $2; }
388 ;

```

- Reduce by $\{3\} \rightarrow \epsilon$, adding the *itype* to the symbol table.
- Shift the SEMI, and read the lookahead symbol. *itype* is in the symbol table this time, so the scanner returns a **TTYPE** token.
- Reduce by *ext_def* $\rightarrow \{ \text{opt_specifiers } \text{ext_decl_list} \} \{3\}$ SEMI.

Listing 6.47. *c.y*— High-Level, External Definitions (Part One)

```

389 ext_def : opt_specifiers ext_decl_list
390     {
391         add_spec_to_decl( $1, $2 );
392
393         if( !$1->tdef )
394             discard_link_chain( $1 );
395
396         add_symbols_to_table ( $2 = reverse_links( $2 ) );
397         figure_osclass ( $2 );
398         generate_defs_and_free_args ( $2 );
399         remove_duplicates ( $2 );
400     }
401     SEMI
402
403     /* There are additional right-hand sides listed in subsequent listings.
404     */

```

The action on lines 390 to 401 of Listing 6.47 needs some discussion. The attribute associated with the *ext_decl_list* at \$2 is a pointer to a linked list of symbol structures, one for each variable in the declarator list. The attribute associated with *opt_specifiers* at \$1 is one of two things: either the specifier component of a declaration, or, if the declaration used a synthetic type, the complete type chain as was stored in the symbol-table entry for the **typedef**. In both cases, the *add_spec_to_decl()* call on line 391 modifies every type chain in the list of symbols by adding a copy of the type chain passed in as the first argument to the end of each symbol's type chain. Then, if the current specifier didn't come from a **typedef**, the extra copy is discarded on line 394. The symbols are added to symbol table on line 396. The *figure_osclass()* call on

line 397 determines the output storage class of all symbols in the chain, generate_defs_and_free_args() outputs the actual C-code definitions, and remove_duplicates() destroys any duplicate declarations in case a declaration and definition of a global variable are both present. All of these subroutines are in Listings 6.48 and 6.49, below.

Listing 6.48. *decl.c*—Symbol-Table Manipulation and C-code Declarations

```

154 void      add_symbols_to_table( sym )
155 symbol   *sym;
156 {
157     /* Add declarations to the symbol table.
158     *
159     * Serious redefinitions (two publics, for example) generate an error
160     * message. Harmless redefinitions are processed silently. Bad code is
161     * generated when an error message is printed. The symbol table is modified
162     * in the case of a harmless duplicate to reflect the higher precedence
163     * storage class: (public == private) > common > extern.
164     *
165     * The sym->rname field is modified as if this were a global variable (an
166     * underscore is inserted in front of the name). You should add the symbol
167     * chains to the table before modifying this field to hold stack offsets
168     * in the case of local variables.
169     */
170
171     symbol *exists;           /* Existing symbol if there's a conflict.      */
172     int    harmless;
173     symbol *new;
174
175     for(new = sym; new ; new = new->next )
176     {
177         exists = (symbol *) findsym(Symbol_tab, new->name);
178
179         if( !exists || exists->level != new->level )
180         {
181             sprintf ( new->rname, "_%1.*s", sizeof(new->rname)-2, new->name );
182             addsym ( Symbol_tab, new );
183         }
184         else
185         {
186             harmless      = 0;
187             new->duplicate = 1;
188
189             if( the_same_type( exists->type, new->type, 0 ) )
190             {
191                 if( exists->etype->OCLASS==EXT || exists->etype->OCLASS==COM )
192                 {
193                     harmless = 1;
194
195                     if( new->etype->OCLASS != EXT )
196                     {
197                         exists->etype->OCLASS = new->etype->OCLASS;
198                         exists->etype->SCLASS = new->etype->SCLASS;
199                         exists->etype->EXTERN = new->etype->EXTERN;
200                         exists->etype->STATIC = new->etype->STATIC;
201                     }
202                 }
203             }
204         }
205     }

```



Listing 6.48. continued...

```

204             if( !harmless )
205                 yyerror("Duplicate declaration of %s\n", new->name );
206         }
207     }
208 }
209
210 /*-----*/
211
212 void      figure_osclass( sym )
213 symbol    *sym;
214 {
215     /* Go through the list figuring the output storage class of all variables.
216      * Note that if something is a variable, then the args, if any, are a list
217      * of initializers. I'm assuming that the sym has been initialized to zeros;
218      * at least the OSCLASS field remains unchanged for nonautomatic local
219      * variables, and a value of zero there indicates a nonexistent output class.
220      */
221
222     for( ; sym ; sym = sym->next )
223     {
224         if( sym->level == 0 )
225         {
226             if( IS_FUNCT( sym->type ) )
227             {
228                 if      ( sym->etype->EXTERN ) sym->etype->OCLASS = EXT;
229                 else if ( sym->etype->STATIC ) sym->etype->OCLASS = PRI;
230                 else                                sym->etype->OCLASS = PUB;
231             }
232             else
233             {
234                 if      ( sym->etype->STATIC ) sym->etype->OCLASS = PRI;
235                 if      ( sym->args           ) sym->etype->OCLASS = PUB;
236                 else                                sym->etype->OCLASS = COM;
237             }
238         }
239         else if( sym->type->SCLASS == FIXED )
240         {
241             if      ( IS_FUNCT ( sym->type ) ) sym->etype->OCLASS = EXT;
242             else if (! IS_LABEL ( sym->type ) ) sym->etype->OCLASS = PRI;
243         }
244     }
245 }
246
247 /*-----*/
248
249 void      generate_defs_and_free_args( sym )
250 symbol    *sym;
251 {
252     /* Generate global-variable definitions, including any necessary
253      * initializers. Free the memory used for the initializer (if a variable)
254      * or argument list (if a function).
255      */
256
257     for( ; sym ; sym = sym->next )
258     {
259         if( IS_FUNCT(sym->type) )
260         {
261             /* Print a definition for the function and discard arguments
262              * (you'd keep them if prototypes were supported).
263              */

```



Listing 6.48. continued...

```

264             yydata( "external\t%s();\n", sym->rname );
265             discard_symbol_chain( sym->args );
266             sym->args = NULL;
267         }
268     }
269     else if( IS_CONSTANT(sym->etype) || sym->type->tdef )
270     {
271         continue;
272     }
273     else if( !sym->args )           /* It's an uninitialized variable. */
274     {
275         print_bss_dcl( sym );       /* Print the declaration. */
276     }
277     else                         /* Deal with an initializer. */
278     {
279         var_dcl( yydata, sym->etype->OCLASS, sym, "=" );
280
281         if( IS_AGGREGATE( sym->type ) )
282             yyerror("Initialization of aggregate types not supported\n");
283
284         else if( !IS_CONSTANT( ((value *)sym->args)->etype ) )
285             yyerror("Initializer must be a constant expression\n");
286
287         else if( !the_same_type(sym->type, ((value *) sym->args)->type, 0) )
288             yyerror("Initializer: type mismatch\n");
289
290         else
291             yydata( "%s;\n", CONST_STR( (value *) sym->args ) );
292
293         discard_value( (value *) (sym->args) );
294         sym->args = NULL;
295     }
296 }
297 }
298
299 /*-----*/
300
301 symbol *remove_duplicates( sym )
302 symbol *sym;
303 {
304     /* Remove all nodes marked as duplicates from the linked list and free the
305      * memory. These nodes should not be in the symbol table. Return the new
306      * head-of-list pointer (the first symbol may have been deleted).
307      */
308
309     symbol *prev = NULL;
310     symbol *first = sym;
311
312     while( sym )
313     {
314         if( !sym->duplicate )           /* Not a duplicate, go to the      */
315         {                            /* next list element.          */
316             prev = sym;
317             sym = sym->next;
318         }
319         else if( prev == NULL )        /* Node is at start of the list. */
320         {
321             first = sym->next;
322             discard_symbol( sym );
323             sym = first;

```



Listing 6.48. continued...

```

324         }
325     else                                /* Node is in middle of the list. */
326     {
327         prev->next = sym->next;
328         discard_symbol( sym );
329         sym = prev->next;
330     }
331 }
332 return first;
333 }
```

Listing 6.49. decl.c—Generate C-Code Definitions

```

334 void      print_bss_dcl( sym )      /* Print a declaration to the bss segment. */
335 symbol    *sym;
336 {
337     if( sym->etype->SCLASS != FIXED )
338     yyerror( "Illegal storage class for %s", sym->name );
339     else
340     {
341         if( sym->etype->STATIC && sym->etype->EXTERN )
342             yyerror(">%s: Bad storage class\n", sym->name );
343
344         var_dcl( yybss, sym->etype->OCLASS, sym, ";" );
345     }
346 }
347
348 /*-----*/
349
350 PUBLIC void var_dcl( ofunct, c_code_sclass, sym, terminator )
351
352 void      (* ofunct)();      /* Pointer to output function (yybss or yydata). */
353 int       c_code_sclass;   /* C-code storage class of symbol. */
354 symbol    *sym;           /* Symbol itself. */
355 char     *terminator;    /* Print this string at end of the declaration. */
356 {
357     /* Kick out a variable declaration for the current symbol. */
358
359     char suffix[32];
360     char *type            = "";
361     int size              = 1 ;
362     link *p                = sym->type;
363     char *storage_class  = (c_code_sclass == PUB) ? "public" : :
364                                         (c_code_sclass == PRI) ? "private" : :
365                                         (c_code_sclass == EXT) ? "external" : "common" ;
366     *suffix = '\0';
367
368     if( IS_FUNCT(p) )
369     {
370         yyerror("INTERNAL, var_dcl: object not a variable\n");
371         exit( 1 );
372     }
373
374     if( IS_ARRAY(p) )
375     {
```

Listing 6.49. continued...

```

376         for(; IS_ARRAY(p) ; p = p->next )
377             size *= p->NUM_ELE ;
378
379         sprintf( suffix, "[%d]", size );
380     }
381
382     if( IS_STRUCT(p) )
383     {
384         (*ofunct )( "\nALIGN(lword)\n" );
385         sprintf( suffix, "[%d]", size * p->V_STRUCT->size );
386     }
387
388     if( IS_POINTER(p) )
389         type = PTYPE;
390     else                                     /* Must be a specifier. */
391         switch( p->NOUN )
392         {
393             case CHAR:      type = CTYPE;           break;
394             case INT:       type = p->LONG ? LTYPE : ITYPE; break;
395             case STRUCTURE: type = STYPE;           break;
396         }
397
398     (*ofunct )( "%s\t%s\t%s%s%s", storage_class, type,
399                           sym->rname, suffix, terminator );
400 }
```

6.6.2 Structure and Union Declarations

Now, let's back up for a moment and look at more-complex types: structures and unions. These were handled at a high level on line 231 of Listing 6.39 on page 525, where a *struct_specifier* is recognized in place of a **TYPE** token. For example, everything from the **struct** up to the name is a *struct_specifier* in the following declaration:

```

struct tag
{
    int tinker;
    long tailor;
    char soldier;
    struct tag *spy;
}
name;
```

A parse for this declaration is shown in Table 6.14. There are several points of interest here. First, the **STRUCT** token is returned by the scanner for both the **struct** and **union** lexemes. The associated `<ascii>` attribute is the first character of the lexeme. In practice, the only difference between a structure and a union is the way that the offsets to each field are computed—in a union the offsets are all zero. Since the two data types are syntactically identical, you need only one token with an attribute that allows the code-generation action to distinguish the two cases.

The structure tag is handled in the parse on lines 4 to 6 of Table 6.14. The associated actions are on lines 426 to 444 of Listing 6.50, below. There are two actions of interest, $opt_tag \rightarrow \epsilon$ on line 427 takes care of those situations where an explicit tag is not present. An arbitrary name is generated on line 430, and a new *struct_def* is allocated on the next line. Then the definition (with no fields as of yet) is added to the structure table on line 433. A pointer to the new structure-table element is passed up as an attribute.

STRUCT recognized for
both **struct** and **union**.
`<ascii>` attribute.

Structure tags.

$opt_tag \rightarrow \epsilon$

Table 6.14. Parsing a Structure Definition

	Stack	Next Action
1	(empty)	Reduce: <i>ext_def_list</i> → ε DATA: #include <tools/virtual.h> DATA: #define T(x) DATA: SEG(data) CODE: SEG(code) BSS: SEG(bss)
2	<i>ext_def_list</i>	Shift: STRUCT
3	<i>ext_def_list</i> STRUCT	Shift: NAME
4	<i>ext_def_list</i> STRUCT NAME	Reduce: tag → NAME
5	<i>ext_def_list</i> STRUCT tag	Reduce: opt_tag → tag
6	<i>ext_def_list</i> STRUCT opt_tag	Shift: LC
7	<i>ext_def_list</i> STRUCT opt_tag LC	Reduce: def_list → ε
8	<i>ext_def_list</i> STRUCT opt_tag LC def_list	Shift: TYPE
9	<i>ext_def_list</i> STRUCT opt_tag LC def_list TYPE	Reduce: type_specifier → TYPE
10	<i>ext_def_list</i> STRUCT opt_tag LC def_list type_specifier	Reduce: type_or_class → type_specifier
11	<i>ext_def_list</i> STRUCT opt_tag LC def_list type_or_class	Reduce: specifiers → type_or_class
12	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers	Shift: NAME
13	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers NAME	Reduce: new_name → NAME
14	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers new_name	Reduce: var_decl → new_name
15	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers var_decl	Reduce: decl → var_decl
16	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl	Reduce: decl_list → decl
17	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list	Reduce: {65} → ε
18	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list {65}	Shift: SEMI
19	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list {65} SEMI	Reduce: def → specifiers decl_list {65} SEMI
20	<i>ext_def_list</i> STRUCT opt_tag LC def_list def	Reduce: def_list → def_list def
21	<i>ext_def_list</i> STRUCT opt_tag LC def_list	Shift: TYPE
22	<i>ext_def_list</i> STRUCT opt_tag LC def_list TYPE	Reduce: type_specifier → TYPE
23	<i>ext_def_list</i> STRUCT opt_tag LC def_list type_specifier	Reduce: type_or_class → type_specifier
24	<i>ext_def_list</i> STRUCT opt_tag LC def_list type_or_class	Reduce: specifiers → type_or_class
25	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers	Shift: NAME
26	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers NAME	Reduce: new_name → NAME
27	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers new_name	Reduce: var_decl → new_name
28	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers var_decl	Reduce: decl → var_decl
29	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl	Reduce: decl_list → decl
30	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list	Reduce: {65} → ε
31	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list {65}	Shift: SEMI
32	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list {65} SEMI	Reduce: def → specifiers decl_list {65} SEMI
33	<i>ext_def_list</i> STRUCT opt_tag LC def_list def	Reduce: def_list → def_list def
34	<i>ext_def_list</i> STRUCT opt_tag LC def_list	Shift: TYPE
35	<i>ext_def_list</i> STRUCT opt_tag LC def_list TYPE	Reduce: type_specifier → TYPE
36	<i>ext_def_list</i> STRUCT opt_tag LC def_list type_specifier	Reduce: type_or_class → type_specifier
37	<i>ext_def_list</i> STRUCT opt_tag LC def_list type_or_class	Reduce: specifiers → type_or_class
38	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers	Shift: NAME
39	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers NAME	Reduce: new_name → NAME
40	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers new_name	Reduce: var_decl → new_name
41	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers var_decl	Reduce: decl → var_decl
42	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl	Reduce: decl_list → decl
43	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list	Reduce: {65} → ε
44	<i>ext_def_list</i> STRUCT opt_tag LC def_list specifiers decl_list {65}	Shift: SEMI

continued...

The action on lines 437 to 444 is used when a tag is present in the definition, as is the case in the current example. The incoming attribute for the NAME token is useless here because this attribute is generated by a symbol-table lookup, not a structure-table lookup. The action code looks up the name in the structure table and returns a pointer to the entry if it's there. Otherwise, a new structure-table element is created and added to

Table 6.14. Continued. Parsing a Structure Definition

	Stack	Next Action
45	<i>ext_def_list STRUCT opt_tag LC def_list specifiers decl_list {65} SEMI</i>	Reduce: $\text{def} \rightarrow \text{specifiers decl_list } \{65\} \text{ SEMI}$
46	<i>ext_def_list STRUCT opt_tag LC def_list def</i>	Reduce: $\text{def_list} \rightarrow \text{def_list def}$
47	<i>ext_def_list STRUCT opt_tag LC def_list</i>	Shift: STRUCT
48	<i>ext_def_list STRUCT opt_tag LC def_list STRUCT</i>	Shift: NAME
49	<i>ext_def_list STRUCT opt_tag LC def_list STRUCT NAME</i>	Reduce: $\text{tag} \rightarrow \text{NAME}$
50	<i>ext_def_list STRUCT opt_tag LC def_list STRUCT tag</i>	Reduce: $\text{struct_specifier} \rightarrow \text{STRUCT tag}$
51	<i>ext_def_list STRUCT opt_tag LC def_list struct_specifier</i>	Reduce: $\text{type_specifier} \rightarrow \text{struct_specifier}$
52	<i>ext_def_list STRUCT opt_tag LC def_list type_specifier</i>	Reduce: $\text{type_or_class} \rightarrow \text{type_specifier}$
53	<i>ext_def_list STRUCT opt_tag LC def_list type_or_class</i>	Reduce: $\text{specifiers} \rightarrow \text{type_or_class}$
54	<i>ext_def_list STRUCT opt_tag LC def_list specifiers</i>	Shift: STAR
55	<i>ext_def_list STRUCT opt_tag LC def_list specifiers STAR</i>	Shift: NAME
56	<i>ext_def_list STRUCT opt_tag LC def_list specifiers STAR NAME</i>	Reduce: $\text{new_name} \rightarrow \text{NAME}$
57	<i>ext_def_list STRUCT opt_tag LC def_list specifiers STAR new_name</i>	Reduce: $\text{var_decl} \rightarrow \text{new_name}$
58	<i>ext_def_list STRUCT opt_tag LC def_list specifiers STAR var_decl</i>	Reduce: $\text{var_decl} \rightarrow \text{STAR var_decl}$
59	<i>ext_def_list STRUCT opt_tag LC def_list specifiers var_decl</i>	Reduce: $\text{decl} \rightarrow \text{var_decl}$
60	<i>ext_def_list STRUCT opt_tag LC def_list specifiers decl</i>	Reduce: $\text{decl_list} \rightarrow \text{decl}$
61	<i>ext_def_list STRUCT opt_tag LC def_list specifiers decl_list</i>	Reduce: $\{65\} \rightarrow \epsilon$
62	<i>ext_def_list STRUCT opt_tag LC def_list specifiers decl_list {65}</i>	Shift: SEMI
63	<i>ext_def_list STRUCT opt_tag LC def_list specifiers decl_list {65} SEMI</i>	Reduce: $\text{def} \rightarrow \text{specifiers decl_list } \{65\} \text{ SEMI}$
64	<i>ext_def_list STRUCT opt_tag LC def_list def</i>	Reduce: $\text{def_list} \rightarrow \text{def_list def}$
65	<i>ext_def_list STRUCT opt_tag LC def_list</i>	Shift: RC
66	<i>ext_def_list STRUCT opt_tag LC def_list RC</i>	Reduce: $\text{struct_specifier} \rightarrow \text{STRUCT opt_tag LC def_list RC}$
67	<i>ext_def_list struct_specifier</i>	Reduce: $\text{type_specifier} \rightarrow \text{struct_specifier}$
68	<i>ext_def_list type_specifier</i>	Reduce: $\text{type_or_class} \rightarrow \text{type_specifier}$
69	<i>ext_def_list type_or_class</i>	Reduce: $\text{specifiers} \rightarrow \text{type_or_class}$
70	<i>ext_def_list specifiers</i>	Reduce: $\text{opt_specifiers} \rightarrow \text{specifiers}$
71	<i>ext_def_list opt_specifiers</i>	Shift: NAME
72	<i>ext_def_list opt_specifiers NAME</i>	Reduce: $\text{new_name} \rightarrow \text{NAME}$
73	<i>ext_def_list opt_specifiers new_name</i>	Reduce: $\text{var_decl} \rightarrow \text{new_name}$
74	<i>ext_def_list opt_specifiers var_decl</i>	Reduce: $\text{ext_decl} \rightarrow \text{var_decl}$
75	<i>ext_def_list opt_specifiers ext_decl</i>	Reduce: $\text{ext_decl_list} \rightarrow \text{ext_decl}$
76	<i>ext_def_list opt_specifiers ext_decl_list</i>	Reduce: $\{50\} \rightarrow \epsilon$ BSS: ALIGN(1word) BSS: common record _name[16];
77	<i>ext_def_list opt_specifiers ext_decl_list {50}</i>	Shift: SEMI
78	<i>ext_def_list opt_specifiers ext_decl_list {50} SEMI</i>	Reduce: $\text{ext_def} \rightarrow \text{opt_specifiers ext_decl_list } \{50\} \text{ SEMI}$
79	<i>ext_def_list ext_def</i>	Reduce: $\text{ext_def_list} \rightarrow \text{ext_def_list ext_def}$
80	<i>ext_def_list</i>	Reduce: $\text{program} \rightarrow \text{ext_def_list}$ (Accept)

the table, as before. This time the real tag is used rather than an arbitrary name, however.

The parser now moves on to process the field definitions, using the productions in Listing 6.51, below. These productions are effectively the same as the *ext_def_list* productions that handle global-variable definitions. In fact, most of the actions are identical. You need two sets of productions because function definitions (with bodies) are permitted only at the global level. As with an *ext_def*, the attribute associated with a *def* is a pointer to a cross-linked chain of symbols, one for each comma-separated declarator found in a definition, or NULL if no declarators were found. The *def_list* productions process a list of semicolon-terminated definitions, linking the cross-linked symbols

Structure-field definitions.

def_list

Listing 6.50. c.y— Structures

```

405 struct_specifier
406     : STRUCT opt_tag LC def_list RC
407     {
408         if( !$2->fields )
409         {
410             $2->fields = reverse_links( $4 );
411
412             if( !illegal_struct_def( $2, $4 ) )
413                 $2->size = figure_struct_offsets( $2->fields, $1=='s' );
414         }
415     else
416     {
417         yyerror("Ignoring redefinition of %s", $2->tag );
418         discard_symbol_chain( $4 );
419     }
420
421     $$ = $2;
422     }
423     | STRUCT tag           { $$ = $2; }
424     ;
425
426 opt_tag : tag
427     | /* empty */    {
428         static unsigned label = 0;
429         static char tag[16];
430         sprintf( tag, "%03d", label++ );
431
432         $$ = new_structdef( tag );
433         addsym( Struct_tab, $$ );
434     }
435     ;
436
437 tag      : NAME      {
438     if( !($$=(structdef *)) findsym(Struct_tab,yytext) )
439     {
440         $$      = new_structdef( yytext );
441         $$->level = Nest_level;
442         addsym( Struct_tab, $$ );
443     }
444     }
445     ;

```

from each individual definition together into a larger list. The new elements are added to the end of the list—the loop on line 450 finds the end. After the entire list is processed (on line 65 of the parse), the attribute associated with the *def_list* is a linked list of symbols, one for each field in the structure.

The parser now reduces by

struct_specifier → **STRUCT** *opt_tag* **LC** *def_list* **RC**

struct_specifier →
STRUCT *opt_tag* **LC**
def_list **RC**

Figuring offsets to fields.

The action is on lines 407 to 421 of Listing 6.50, and the subroutines that are used here are in Listing 6.52. The *illegal_struct_def()* call checks the field definitions to make sure that there's no recursion and that none of the fields are function definitions (as compared to function pointers, which are legal). The *figure_struct_offsets()* call on line 427 figures the offsets from the base address of the structure to the individual fields. The basic algorithm just traverses the linked list of symbols adding the

Listing 6.51. c.y—Local Variables and Function Arguments

```

446 def_list
447     : def_list def          { symbol *p;
448                               if( p = $2 )
449                               {
450                                   for( p->next; p = p->next )
451                                       ;
452                                   p->next = $1;
453                                   $$      = $2;
454                               }
455                           }
456     | /* epsilon */        { $$ = NULL; } /* Initialize end-of-list */
457     ;                      /* pointer. */
458
459 def
460     : specifiers decl_list { add_spec_to_decl( $1, $2 );   }
461           SEMI { $$ = $2; }                         }
462
463     | specifiers SEMI     { $$ = NULL; }
464     ;
465
466 decl_list
467     : decl                  { $$->next = NULL; }
468     | decl_list COMMA decl
469     {
470         $3->next = $1;
471         $$      = $3;
472     }
473     ;
474
475 decl
476     : funct_decl
477     | var_decl
478     | var_decl EQUAL initializer { yyerror( "Ignoring initializer.\n");
479                                     discard_value( $3 );
480                                     }
481     | var_decl COLON const_expr    %prec COMMA
482     | COLON const_expr            %prec COMMA
483     ;

```

cumulative size of the preceding fields to the current offset. A minor problem is caused by alignment restrictions—padding may have to be added in the middle of the structure in order to get interior fields aligned properly. The structure is declared as a character array on line 76 of the parse (record is an alias for byte), and the individual fields are not declared explicitly—they are extracted from the proper location within the array when expressions are processed. As a consequence, the compiler has to worry about supplying padding that would normally be supplied by the assembler.

This alignment problem actually arises in the current example. The first field is a two-byte int, but the second field requires four-byte alignment, so two spaces of padding have to be added to get the second field aligned properly. The situation is simplified, somewhat, by assuming that the first field is always aligned on a worst-case boundary. An align(1word) directive is generated by the compiler just above the actual variable definition for this purpose. Finally, note that the structure size is rounded up to be an even multiple of the worst-case alignment restriction on lines 464 and 465 of Listing 6.52 so that arrays of structures work correctly.

Structure-field alignment, padding.

Structure size rounded up for arrays.

optSpecifier,

The structure definition is now reduced to an *optSpecifier*, and the parse continues just like a simple variable definition. One more right-hand side to *extDef* is needed for structures. It is shown in Listing 6.53, and handles structure, union, and enumerated-type declarations that don't allocate space for a variable (such as a **struct** definition with a tag but no variable name). Note that the *discardLinkChain()* call on line 490 does not delete anything from the structure table.

Listing 6.52. *decl.c*—Structure-Processing Subroutines

```

401 int illegal_struct_def( cur_struct, fields )
402 structdef *cur_struct;
403 symbol *fields;
404 {
405     /* Return true if any of the fields are defined recursively or if a function
406      * definition (as compared to a function pointer) is found as a field.
407      */
408
409     for( ; fields; fields = fields->next )
410     {
411         if( IS_FUNCT(fields->type) )
412         {
413             yyerror("struct/union member may not be a function");
414             return 1;
415         }
416         if( IS_STRUCT(fields->type) &&
417             !strcmp( fields->type->V_STRUCT->tag, cur_struct->tag ) )
418         {
419             yyerror("Recursive struct/union definition\n");
420             return 1;
421         }
422     }
423     return 0;
424 }
425 -----
426
427 int figure_struct_offsets( p, is_struct )
428 symbol *p;                                /* Chain of symbols for fields. */
429 int is_struct;                            /* 0 if a union. */
430 {
431     /* Figure the field offsets and return the total structure size. Assume
432      * that the first element of the structure is aligned on a worst-case
433      * boundary. The returned size is always an even multiple of the worst-case
434      * alignment. The offset to each field is put into the "level" field of the
435      * associated symbol.
436      */
437
438     int align_size, obj_size;
439     int offset = 0;
440
441     for( ; p ; p = p->next )
442     {
443         if( !is_struct )                      /* It's a union. */
444         {
445             offset = max( offset, get_sizeof( p->type ) );
446             p->level = 0;
447         }
448     else
449     {
450         obj_size = get_sizeof( p->type );
451         align_size = get_alignment( p->type );

```



Listing 6.52. continued...

```

452
453         while( offset % align_size )
454             ++offset;
455
456         p->level  = offset;
457         offset   += obj_size ;
458     }
459 }
460 /* Return the structure size: the current offset rounded up to the      */
461 /* worst-case alignment boundary. You need to waste space here in case  */
462 /* this is an array of structures.                                         */
463
464     while( offset % ALIGN_WORST )
465         ++offset ;
466     return offset;
467 }
468 /*-----*/
469
470 int      get_alignment( p )
471 link    *p;
472 {
473     /* Returns the alignment--the number by which the base address of the object
474      * must be an even multiple. This number is the same one that is returned by
475      * get_sizeof(), except for structures which are worst-case aligned, and
476      * arrays, which are aligned according to the type of the first element.
477      */
478
479     int size;
480
481     if( !p )
482     {
483         yyerror("INTERNAL, get_alignment: NULL pointer\n");
484         exit( 1 );
485     }
486     if( IS_ARRAY( p )           ) return get_alignment( p->next );
487     if( IS_STRUCT( p )        ) return ALIGN_WORST;
488     if( size = get_sizeof( p ) ) return size;
489
490     yyerror("INTERNAL, get_alignment: Object aligned on zero boundary\n");
491     exit( 1 );
492 }
```

Listing 6.53. c.y— High-Level, External Definitions (Part Two)

```

484 /* ext_def :                                                 */
485     | opt_specifiers
486     {
487         if( !($1->class == SPECIFIER && $1->NOUN == STRUCTURE) )
488             yyerror("Useless definition (no identifier)\n");
489         if( !$1->tdef )
490             discard_link_chain( $1 );
491     }
492 SEMI
```

6.6.3 Enumerated-Type Declarations

The final definition recognized by the grammar is an enumerated type, handled at a high level on line 230 of Listing 6.39 (on page 525) where an *enum_specifier* is recognized in place of a **TYPE** token. An enumerated-type definition like this:

```
enum tag { rich_man, poor_man, beggar_man = 5, thief } x;
```

is treated as if the following had been used:

```
int x;
#define rich_man      0
#define poor_man     (rich_man + 1)
#define beggar_man   5
#define thief        (beggar_man + 1)
```

but the compiler recognizes the elements of the enumerated type directly rather than using a macro preprocessor. The high-level action in Listing 6.39 on page 525 just creates a specifier for an **int**, ignoring the tag component of the enumerated type.

Enumerated types are handled as integer constants.

`Enum_val, do_enum().`

The real work happens in the productions in Listing 6.54, which creates symbol-table entries for the symbolic constants (`rich_man`, and so forth). Internally, the compiler doesn't distinguish between an element of an enumerated type and any other integer constant. When an enumerated-type element is referenced, a symbol-table lookup is necessary to get the value; but thereafter, the value is handled just like any other integer constant. The `Enum_val` global variable keeps track of the current constant value. It is initialized to zero on line 507 of Listing 6.54, when the **enum** keyword is recognized. The enumerated-type elements are processed on lines 516 to 518. `Enum_val` is modified as necessary if an explicit value is given in the definition. In any event, `do_enum()`, in Listing 6.56, is called to create the symbol-table element, which is added to the table on the next line.

Listing 6.54. *c.y*— Enumerated Types

```
122  %{
123  int      Enum_val;           /* Current enumeration constant value */
124  %}

493 enum_specifier
494     : enum name opt_enum_list    { if( $2->type )
495                                         yyerror("%s: redefinition", $2->name);
496                                         else
497                                             discard_symbol($2);
498                                         }
499     | enum LC enumerator_list RC
500     ;
501
502 opt_enum_list
503     : LC enumerator_list RC
504     | /* empty */
505     ;
506
507 enum     : ENUM    { Enum_val = 0; }
508     ;
509
510 enumerator_list
511     : enumerator
512     | enumerator_list COMMA enumerator
513     ;
```



Listing 6.54. continued...

```

515 enumerator
516     : name           { do_enum( $1, Enum_val++ ); }
517     | name EQUAL const_expr { Enum_val = $3;
518                               do_enum( $1, Enum_val++ ); }
519 ;

```

Listing 6.55. decl.c— Enumerated-Type Subroutines

```

493 PUBLIC void do_enum( sym, val )
494 symbol *sym;
495 int val;
496 {
497     if( conv_sym_to_int_const( sym, val ) )
498         addsym( Symbol_tab, sym );
499     else
500     {
501         yyerror( "%s: redefinition", sym->name );
502         discard_symbol( sym );
503     }
504 }
505 /*-----*/
506 PUBLIC int conv_sym_to_int_const( sym, val )
507 symbol *sym;
508 int val;
509 {
510     /* Turn an empty symbol into an integer constant by adding a type chain
511      * and initializing the v_int field to val. Any existing type chain is
512      * destroyed. If a type chain is already in place, return 0 and do
513      * nothing, otherwise return 1. This function processes enum's.
514 */
515     link *lp;
516
517     if( sym->type )
518         return 0;
519     lp = new_link();
520     lp->class = SPECIFIER;
521     lp->NOUN = INT;
522     lp->SCLASS = CONSTANT;
523     lp->V_INT = val ;
524     sym->type = lp;
525     *sym->rname = '\0';
526     return 1;
527 }

```

6.6.4 Function Declarations

The next sort of declaration is a function declaration, handled, at the high level, by the remaining right-hand side to *ext_def*:

```
ext_def → opt_specifiers funct_decl { . . . } def_list { . . . } compound_stmt { . . . };
```

We'll look at the actual production in a moment. First, look at the parse of the following code in Table 6.15:

```
pooh(piglet, eeyore)
long eeyore;
{  
}
```

This time I've shown the output as well as the parser actions; the three output streams are indicated by **CODE:**, **DATA:**, and **BSS:** in the table. The final compiler output (after the three streams have been merged) is in Listing 6.56.

Labels for `link`: generated after processing subroutine, printed above subroutine.

Output streams: `code`, `data`, `bss`.

Subroutine-prefix-and-suffix generation.

Prefix.

Suffix.

Arguments to `link` instruction.

`Vspace`, `Tspace`, `Funct_name`.

ext_def → *opt_specifiers funct_decl {} def_list {} compound_stmt {}.*

Note that the time at which the parser outputs something is not directly related to the position of that something in the output file. For example, the **#defines** for `L0` and `L1` at the top of the output file are not emitted by the parser until after the entire subroutine has been processed. This shuffling is accomplished by using two different output streams. The definitions are written to the `bss` stream; code is written to the `code` stream. The streams are merged in the following order: `data`, `bss`, and then `code`. So, all output to the `bss` stream appears above all `code`-stream output in the final program.

The code on lines nine to 14 of Listing 6.56 is generated at the top of every subroutine (with the obvious customizations). This block of instructions is called the subroutine *prefix*. The macro definitions on lines five and six of Listing 6.56 are also part of the prefix, even though these definitions end up at the top of the output file rather than immediately above the subroutine definition, because they are generated along with the rest of the prefix code. The code on lines 18 to 20 of Listing 6.56 is generated at the bottom of every subroutine and is called the subroutine *suffix*.

The `L0` and `L1` macros on lines five and six of Listing 6.56 are used as arguments in the `link` instruction at the top of the subroutine. The numeric component of the labels is unique—the next subroutine in the input file uses `L2` and `L3`. In the current example, `L0` holds the size, in bytes, of the local-variable region of the subroutine's stack frame; `L1` holds the size of the temporary-variable region. Unfortunately, for reasons that I'll discuss in a moment, neither size is known until after the entire subroutine has been processed. The compiler solves the problems by generating the `L0` and `L1` labels at the same time that it outputs the `link` instruction. The label names are stored internally in the `Vspace` and `Tspace` arrays in Listing 6.57. The compiler puts the function name into `Funct_name` at the same time. Later on in the parse, after the subroutine has been processed and the sizes of the two regions are known, the compiler emits the **#defines** to the `data` segment, using the previously generated labels.¹⁷

The code to do all of the foregoing is in the first action of the following production:

```
ext_def → opt_specifiers funct_decl { . . . } def_list { . . . } compound_stmt { . . . };
```

found on lines 522 to 546 of Listing 6.58. The attribute attached to the *funct_decl* at \$2

17. This technique is necessary only if you are generating code for a single-pass assembler. Two-pass assemblers let you define a label after it's used—the definition is picked up in the first pass and substitutions are made in the second. The two-pass VAX/UNIX assembler just generates the labels at the end of the subroutine code.

Table 6.15. A Parse of pooh(piglet, eeyore) **long** eeyore; {}

	Stack	Next Action
1	(empty)	Reduce: <i>ext_def_list</i> → ϵ DATA: #include <tools/virtual.h> DATA: #define T(x) DATA: SEG(data) CODE: SEG(code) BSS: SEG(bss)
2	<i>ext_def_list</i>	Reduce: <i>opt_specifiers</i> → ϵ
3	<i>ext_def_list opt_specifiers</i>	Shift: NAME
4	<i>ext_def_list opt_specifiers NAME</i>	Reduce: <i>new_name</i> → NAME
5	<i>ext_def_list opt_specifiers new_name</i>	Shift: LP
6	<i>ext_def_list opt_specifiers new_name LP</i>	Reduce: {28} → ϵ
7	<i>ext_def_list opt_specifiers new_name LP {28}</i>	Shift: NAME
8	<i>ext_def_list opt_specifiers new_name LP {28} NAME</i>	Reduce: <i>new_name</i> → NAME
9	<i>ext_def_list opt_specifiers new_name LP {28} new_name</i>	Reduce: <i>name_list</i> → <i>new_name</i>
10	<i>ext_def_list opt_specifiers new_name LP {28} name_list</i>	Shift: COMMA
11	<i>ext_def_list opt_specifiers new_name LP {28} name_list COMMA</i>	Shift: NAME
12	<i>ext_def_list opt_specifiers new_name LP {28} name_list COMMA NAME</i>	Reduce: <i>new_name</i> → NAME
13	<i>ext_def_list opt_specifiers new_name LP {28} name_list COMMA new_name</i>	Reduce: <i>name_list</i> → <i>name_list COMMA new_name</i>
14	<i>ext_def_list opt_specifiers new_name LP {28} name_list</i>	Reduce: {29} → ϵ
15	<i>ext_def_list opt_specifiers new_name LP {28} name_list {29}</i>	Shift: RP
16	<i>ext_def_list opt_specifiers new_name LP {28} name_list {29} RP</i>	Reduce: <i>funct_decl</i> → <i>new_name LP {28} name_list {29} RP</i>
17	<i>ext_def_list opt_specifiers funct_decl</i>	Reduce: {51} → ϵ CODE: #undef T CODE: #define T(n) (fp-L0-(n*4)) CODE: PROC(_pooh,public) CODE: link(L0+L1);
18	<i>ext_def_list opt_specifiers funct_decl {51}</i>	Reduce: <i>def_list</i> → ϵ
19	<i>ext_def_list opt_specifiers funct_decl {51} def_list</i>	Shift: TYPE
20	<i>ext_def_list opt_specifiers funct_decl {51} def_list TYPE</i>	Reduce: <i>typeSpecifier</i> → TYPE
21	<i>ext_def_list opt_specifiers funct_decl {51} def_list typeSpecifier</i>	Reduce: <i>typeOrClass</i> → <i>typeSpecifier</i>
22	<i>ext_def_list opt_specifiers funct_decl {51} def_list typeOrClass</i>	Reduce: <i>specifiers</i> → <i>typeOrClass</i>
23	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers</i>	Shift: NAME
24	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers NAME</i>	Reduce: <i>new_name</i> → NAME
25	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers new_name</i>	Reduce: <i>varDecl</i> → <i>new_name</i>
26	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers varDecl</i>	Reduce: <i>decl</i> → <i>varDecl</i>
27	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers decl</i>	Reduce: <i>declList</i> → <i>decl</i>
28	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers declList</i>	Reduce: {65} → ϵ
29	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers declList {65}</i>	Shift: SEMI
30	<i>ext_def_list opt_specifiers funct_decl {51} def_list specifiers declList {65} SEMI</i>	Reduce: <i>def</i> → <i>specifiers declList {65} SEMI</i>
31	<i>ext_def_list opt_specifiers funct_decl {51} def_list def</i>	Reduce: <i>defList</i> → <i>defList def</i>
32	<i>ext_def_list opt_specifiers funct_decl {51} def_list</i>	Reduce: {52} → ϵ CODE: /* fp+4 = piglet [arg] */ CODE: /* fp+8 = eeyore [arg] */
33	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52}</i>	Shift: LC
34	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52} LC</i>	Reduce: {71} → ϵ

continued...

Table 6.15. Continued. A Parse of pooh(piglet, eeyore) **long** eeyore; {}

	Stack	Next Action
35	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52} LC {71}</i>	Reduce: <i>def_list</i> → ε
36	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52} LC {71} def_list</i>	Reduce: <i>local_defs</i> → <i>def_list</i>
37	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52} LC {71} local_defs</i>	Reduce: <i>stmt_list</i> → ε
38	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52} LC {71} local_defs stmt_list</i>	Shift: RC
39	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52} LC {71} local_defs stmt_list RC</i>	Reduce: <i>compound_stmt</i> → LC {71} local_defs stmt_list RC
40	<i>ext_def_list opt_specifiers funct_decl {51} def_list {52} compound_stmt</i>	Reduce: <i>ext_def</i> → <i>opt_specifiers funct_decl {51} def_list {52} compound_stmt</i> CODE: unlink(); CODE: ret(); CODE: ENDP (_pooh) BSS: #define L0 0 /* pooh loc. */ BSS: #define L1 0 /* pooh tmp. */
41	<i>ext_def_list ext_def</i>	Reduce: <i>ext_def_list</i> → <i>ext_def_list ext_def</i>
42	<i>ext_def_list</i>	Reduce: <i>program</i> → <i>ext_def_list</i> Accept

Listing 6.56. Compiler Output for Function Definition

```

1 #include <tools/virtual.h>
2 #define T(x)
3 SEG(data)
4 SEG(bss)
5 #define L0 0 /* pooh: locals */
6 #define L1 0 /* pooh: temps. */
7 SEG(code)
8
9 #undef T
10 #define T(n) (fp-L0-(n*4))
11 PROC(_pooh,public)
12     link(L0+L1);
13             /* fp+4 = piglet [argument] */
14             /* fp+8 = eeyore [argument] */
15
16     /* Code from the body of the subroutine goes here. */
17
18     unlink();
19     ret();
20 ENDP (_pooh)

```

`symbol.args`, cross links.

is a pointer to a `symbol`, and the `args` field of that structure is itself a pointer to a linked list of additional `symbol` structures, one for each argument. The arguments, along with an entry for the function itself, are put into the symbol table on lines 530 to 532 of Listing 6.58. Putting the elements in the table does not affect the cross links. The arguments still form a linked list after the insertion. The `gen()` calls on line 540 and 541 actually emit the `PROC` and `link` directives—I'll come back to this subroutine later, when I discuss expression processing. The `L` prefix in the variable- and temporary-space labels, is put into the labels on lines 535 and 536 of Listing 6.58. `L_LINK` is defined along with several other label prefixes in `label.h`, Listing 6.59. These other label prefixes are used

Alphabetic components of output labels, `label.h`.

Listing 6.57. *c.y*— Global Variables for Function-Declaration Processing (from *Occs Definitions Section*)

```

125  %{
126  char   Vspace[16];
127  char   Tspace[16];      /* The compiler doesn't know the stack-frame size
128          * when it creates a link() directive, so it outputs
129          * a link(VSPACE+TSPACE). Later on, it #defines VSPACE
130          * to the size of the local-variable space and TSPACE
131          * to the size of the temporary-variable space. Vspace
132          * holds the actual name of the VSPACE macro, and
133          * Tspace the TSPACE macro. (There's a different name
134          * for each subroutine.)
135          */
136
137  char   Funct_name[ NAME_MAX+1 ];      /* Name of the current function */
138  %}

```

for processing **if** statements, **while** loops, and so on. I'll discuss them further when the flow control statements are presented.

The parser now handles the K&R-style argument definitions. The new symbol-table entries are all of type **int**, because the *funct_decl* contains a *name_list*, not a *var_list*. Since no types are specified in the input argument list, **int** is supplied automatically. The parser processes the *def_list* next. When the *def_list*-processing is finished, the attribute at \$4—the previous action was \$3—will be a pointer to a linked list of symbols, one for each formal definition. The *fix_types_and_discard_syms()* call on line 549 of Listing 6.58 looks up each of the redefined symbols in the symbol table. If the symbol is there, the type is modified to reflect the redefinition; if it's not, an error message is generated. This subroutine also discards the symbol structures used for the redefinitions. The *figure_param_offsets()* call on line 550 traverses the argument list again because some elements of the list have just been modified to have new types. It patches the symbol structures so that the *rname* field holds an expression that can be used in a C-code instruction to access the variable. All these expressions are relative to the frame pointer (*WP(fp-8)* and so forth). The position in the argument list determines the value of the offset. The *print_offset_comment()* on line 551 prints the comment in the middle of the output that shows what all of these offsets are. All three subroutines are in Listing 6.60.

Skipping past the function body and local-variable processing for a moment (all these are done in the *compound_stmt*), the end-of-function processing is all done in the third action in Listing 6.58, on lines 556 to 575. The *remove_symbols_from_table()* call on line 560 deletes all the subroutine arguments from the symbol table. The subroutine itself stays in the table, however—it's a global-level symbol. The *discard_symbol_chain()* call on the next line frees the memory used for the associated symbol structures and associated type chains. The structure-table is not modified, so the structure definition persists even though the variable doesn't. Finally, the end-of-function code is output with the three *gen()* calls on the following lines, and the link-instruction labels are emitted on lines 568 to 571 of Listing 6.58.

K&R-style argument definitions.

End-of-function processing.

link labels emitted.

Listing 6.58. c.y— High-Level, External Definitions (Part Three)

```

520 /* ext_def :                                     */
521     | opt_specifiers funct_decl
522     {
523         static unsigned link_val = 0;      /* Labels used for link args.*/
524
525         add_spec_to_decl( $1, $2 );      /* Merge the specifier and   */
526                                         /* declarator. */
527         if( !$1->tdef )
528             discard_link_chain( $1 );    /* Discard extra specifier. */
529
530         figure_osclass      ( $2 );      /* Update symbol table. */
531         add_symbols_to_table ( $2 );      /* Add function itself. */
532         add_symbols_to_table ( $2->args ); /* Add the arguments. */
533
534         strcpy ( Funct_name, $2->name );
535         sprintf( Vspace, "%s%d", L_LINK, link_val++ );
536         sprintf( Tspace, "%s%d", L_LINK, link_val++ );
537
538         yycode( "\n#undef T\n" );
539         yycode( "#define T(n) (fp-%s-(n*4))\n\n", Vspace );
540         gen( "PROC", $2->rname, $2->etype->STATIC ? "private":"public");
541         gen( "link", Vspace, Tspace );
542
543         ++Nest_lev;          /* Make nesting level of definition_list
544                               * match nesting level in the funct_decl
545                               */
546     }
547     def_list
548     {
549         fix_types_and_discard_syms ( $4 );
550         figure_param_offsets      ( $2->args );
551         print_offset_comment     ( $2->args, "argument" );
552
553         --Nest_lev;           /* It's incremented again in the compound_stmt */
554     }
555     compound_stmt
556     {
557         purge_undecl();        /* Deal with implicit declarations */
558                                         /* and undeclared symbols. */
559
560         remove_symbols_from_table ( $2->args ); /* Delete arguments. */
561         discard_symbol_chain    ( $2->args );
562
563         gen( ":%s%d", L_RET, rlabel(1) );      /* End-of-function */
564         gen( "unlink" );                      /* code. */
565         gen( "ret" );
566         gen( "ENDP", $2->rname );
567
568         yybss ( "\n#define %s %d\t/* %s: locals */\n",
569                 Vspace, loc_var_space(), $2->name );
570         yybss ( "#define %s %d\t/* %s: temps. */\n",
571                 Tspace, tmp_var_space(), $2->name );
572
573         tmp_reset();           /* Reset temporary-variable system.*/
574                                         /* (This is just insurance.) */
575     }
576 ;

```

Listing 6.59. *label.h*—Output-label Definitions

```

1  /* This file contains definitions for the various label prefixes. All labels
2   * take the form: <prefix><number>, the <number> supplied by the code-generation
3   * action. The prefixes are defined here.
4   */
5
6 #define L_BODY          "BDY"    /* Top of the body of a for loop.      */
7 #define L_COND_END       "QE"     /* End of conditional.               */
8 #define L_COND_FALSE     "QF"     /* True part of conditional (?:).   */
9 #define L_DOEXIT         "DXIT"   /* Just after the end of the do/while. */
10 #define L_DOTEST        "DTST"   /* Just above the test in a do/while. */
11 #define L_DOTOP          "DTOP"   /* Top of do/while loop.            */
12 #define L_ELSE           "EL"     /* Used by else processing.        */
13 #define L_END            "E"      /* End of relational/logical op.   */
14 #define L_FALSE          "F"      /* False target of relational/logical op. */
15 #define L_INCREMENT      "INC"    /* Just above the increment part of for loop. */
16 #define L_LINK           "L"      /* Offset passed to link instruction. */
17 #define L_NEXT           "EXIT"   /* Outside of loop, end of if clause. */
18 #define L_RET            "RET"    /* Above clean-up code at end of subroutine. */
19 #define L_STRING          "S"      /* Strings.                         */
20 #define L_SWITCH          "SW"     /* Used for switches.               */
21 #define L_TEST            "TST"    /* Above test in while/for/if.      */
22 #define L_TRUE            "T"      /* True target of relational/logical operator. */
23 #define L_VAR             "V"      /* Local-static variables.          */

```

Listing 6.60. *decl.c*—Process Subroutine Arguments

```

528 void fix_types_and_discard_syms( sym )
529 symbol *sym;
530 {
531     /* Patch up subroutine arguments to match formal declarations.
532     *
533     * Look up each symbol in the list. If it's in the table at the correct
534     * level, replace the type field with the type for the symbol in the list,
535     * then discard the redundant symbol structure. All symbols in the input
536     * list are discarded after they're processed.
537     *
538     * Type checking and automatic promotions are done here, too, as follows:
539     *     chars are converted to int.
540     *     arrays are converted to pointers.
541     *     structures are not permitted.
542     *
543     * All new objects are converted to autos.
544     */
545
546     symbol *existing, *s;
547
548     while( sym )
549     {
550         if( !( existing = (symbol *)findsym( Symbol_tab,sym->name) )
551             || sym->level != existing->level )
552         {
553             yyerror("%s not in argument list\n", sym->name );
554             exit(1);
555         }

```



Listing 6.60. continued...

```

556         else if( !sym->type || !sym->etype )
557     {
558         yyerror("INTERNAL, fix_types: Missing type specification\n");
559         exit(1);
560     }
561     else if( IS_STRUCT(sym->type) )
562     {
563         yyerror("Structure passing not supported, use a pointer\n");
564         exit(1);
565     }
566     else if( !IS_CHAR(sym->type) )
567     {
568         /* The existing symbol is of the default int type, don't redefine
569          * chars because all chars are promoted to int as part of the call,
570          * so can be represented as an int inside the subroutine itself.
571          */
572
573         if( IS_ARRAY(sym->type) )           /* Make it a pointer to the */
574             sym->type->DCL_TYPE = POINTER;   /* first element.           */
575
576         sym->etype->SCLASS = AUTO;        /* Make it an automatic var. */
577
578         discard_link_chain(existing->type); /* Replace existing type      */
579         existing->type    = sym->type;       /* chain with the current one.*/
580         existing->etype    = sym->etype;
581         sym->type = sym->etype = NULL;      /* Must be NULL for discard_ */
582     }                                     /* symbol() call, below.      */
583     s = sym->next;
584     discard_symbol( sym );
585     sym = s;
586 }
587 }
588
589 /*-----*/
590
591 int figure_param_offsets( sym )
592 symbol *sym;
593 {
594     /* Traverse the chain of parameters, figuring the offsets and initializing
595      * the real name (in sym->rname) accordingly. Note that the name chain is
596      * assembled in reverse order, which is what you want here because the
597      * first argument will have been pushed first, and so will have the largest
598      * offset. The stack is 32 bits wide, so every legal type of object will
599      * require only one stack element. This would not be the case were floats
600      * or structure-passing supported. This also takes care of any alignment
601      * difficulties.
602      *
603      * Return the number of 32-bit stack words required for the parameters.
604      */
605
606     int offset = 4;                      /* First parameter is always at BP(fp+4). */
607     int i;
608
609     for(; sym ; sym = sym->next )
610     {
611         if( IS_STRUCT(sym->type) )
612         {
613             yyerror("Structure passing not supported\n");
614             continue;
615         }

```



Listing 6.60. continued...

```

616
617     sprintf( sym->rname, "fp+%d", offset );
618     offset += SWIDTH ;
619 }
620
621 /* Return the offset in stack elements, rounded up if necessary. */
622
623 return( (offset / SWIDTH) + (offset % SWIDTH != 0) );
624 }
625
626 /* -----
627
628 void print_offset_comment( sym, label )
629 symbol *sym;
630 char *label;
631 {
632     /* Print a comment listing all the local variables. */
633
634     for( ; sym ; sym = sym->next )
635         yycode( "\t/* %16s = %-16s [%s] */\n", sym->rname, sym->name, label );
636 }
```

6.6.5 Compound Statements and Local Variables

The next issue is the function body, which consists of a *compound_stmt*. As you can see from Listing 6.61, below, a compound statement is a list of statements (a *stmt_list*) surrounded by curly braces. Local-variable definitions (*local_defs*) can appear at the beginning of any compound statement, and since this same production is also used to process multiple-statement bodies of loops, and so forth, local variables can be defined at any nesting level.¹⁸ The inner variables must shadow the outer ones until the compiler leaves its scoping level, however. The local version of the variable is used instead of another variable declared at a more outer level with an identical name.

It's difficult, though certainly possible, to modify the size of the stack frame every time that the compiler enters or leaves a scoping level. The main difficulty is temporary-variable management, which is much easier to do if the temporary variable space doesn't move around or change size during the life of a subroutine. My solution is to allocate space for all local variables, regardless of the scoping level, with the single link instruction at the top of the subroutine. From a space-allocation perspective, all variables are treated as if they were declared in the outermost scoping level at run time. At compile time, however, the symbol-table entry for an inner variable is created when that variable is declared, and it is deleted when the compiler leaves the scoping level for that variable. Even though an inner variable continues to exist at run time, that variable cannot be accessed from outside the compound statement because the symbol-table entry for the variable won't exist.

The obvious problem with this approach is that memory is wasted. In the following fragment, for example, the stack region used to store *castor* could be recycled for use by *pollux*; it isn't:

Function bodies,
compound_stmt.

Nested variable
definitions.

Handling scoping levels,
memory allocation.

¹⁸. This is quite-legal C, though the practice is discouraged because it makes it difficult to find the variable definitions when they're buried in a subroutine.

Listing 6.61. *c.y*—Compound Statements

```

577 compound_stmt
578     : LC
579             { ++Nest_lev;
580             loc_reset();
581         }
582         local_defs stmt_list RC { --Nest_lev;
583             remove_symbols_from_table ( $3 );
584             discard_symbol_chain ( $3 );
585         }
586     ;
587 local_defs
588     : def_list { add_symbols_to_table ( $$ = reverse_links( $1 ) );
589                 figure_local_offsets ( $$, Funct_name );
590                 create_static_locals ( $$, Funct_name );
591                 print_offset_comment ( $$, "variable" );
592             }
593     ;

```

```

    {
        {
            int castor;
        }
        {
            int pollux;
        }
    }
}

```

Local-variable definitions are handled by the productions in Listing 6.61 and the subroutines in Listing 6.62. The *def_list* nonterminal on line 588 of Listing 6.61 is the same production that's used for structure fields. Its attribute is a pointer to the head of a linked list of symbols, one for each declaration. These symbols are added to the symbol table on line 588 in Listing 6.61.

Listing 6.62. *local.c*—Local-Variable Management

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/debug.h>
4 #include <tools/hash.h>
5 #include <tools/l.h>
6 #include <tools/compiler.h>
7 #include <tools/c-code.h>
8 #include "syntab.h"
9 #include "proto.h"
10 #include "label.h"
11
12 /* LOCAL.C Subroutines in this file take care of local-variable management. */
13
14 PRIVATE int      Offset = 0 ;      /* Offset from the frame pointer (which also */
15                                /* marks the base of the automatic-variable */
16                                /* region of the stack frame) to the most */
17                                /* recently allocated variable. Reset to 0 */
18                                /* by loc_reset() at the head of every sub- */
19                                /* routine. */
20

```



Listing 6.62. continued...

```

21  extern void      yycode(), yydata(), yybss(), yycomment();
22  /*-----*/
23
24  void      loc_reset()
25  {
26      /* Reset everything back to the virgin state. Call this subroutine just */
27      /* before processing the outermost compound statement in a subroutine. */
28
29      Offset = 0 ;
30  }
31  /*-----*/
32
33  int       loc_var_space()
34  {
35      /* Return the total cumulative size of the temporary-variable region in
36      * stack elements (not bytes). This call outputs the value of the macro
37      * that specifies the variable-space size in the link instruction. Calling
38      * loc_reset() also resets the return value of this subroutine to zero.
39      */
40
41      return( (Offset + (SWIDTH-1)) / SWIDTH );
42  }
43  /*-----*/
44
45  void      figure_local_offsets( symbol *sym, char *funct_name )
46  symbol    *sym;
47  char      *funct_name;
48  {
49      /* Add offsets for all local automatic variables in the sym list. */
50
51      for( symbol sym ; sym = sym->next )
52          if( !IS_FUNCT( sym->type ) && !sym->etype->STATIC )
53              loc_auto_create( sym );
54  }
55  /*-----*/
56
57  void      loc_auto_create( symbol *sym )
58  symbol    *sym;
59  {
60      /* Create a local automatic variable, modifying the "rname" field of "sym"
61      * to hold a string that can be used as an operand to reference that
62      * variable. This name is a correctly aligned reference of the form:
63      *
64      *           fp + offset
65      *
66      * Local variables are packed as well as possible into the stack frame,
67      * though, as was the case with structures, some padding may be necessary
68      * to get things aligned properly.
69      */
70
71      int align_size = get_alignment( sym->type );
72
73      Offset += get_sizeof( sym->type );           /* Offset from frame pointer */
74                                         /* to variable. */
75
76      while( Offset % align_size )                 /* Add any necessary padding */
77          ++Offset;                                /* to guarantee alignment. */
78

```



Listing 6.62. continued...

```

79     sprintf( sym->rname, "fp-%d", Offset );           /* Create the name.      */
80     sym->etype->SCLASS = AUTO;
81 }
82 /*-----*/
83
84 void    create_static_locals( sym, funct_name )
85 symbol  *sym;
86 char    *funct_name;
87 {
88     /* Generate definitions for local, static variables in the sym list. */
89
90     for( sym ; sym = sym->next )
91         if( !IS_FUNCT( sym->type ) && sym->etype->STATIC )
92             loc_static_create( sym, funct_name );
93 }
94 /*-----*/
95
96 void    loc_static_create( sym, funct_name )
97 symbol  *sym;
98 char    *funct_name;
99 {
100     static unsigned val;          /* Numeric component of arbitrary name. */
101
102     sprintf( sym->rname, "%s%d", L_VAR, val++ );
103     sym->etype->SCLASS = FIXED ;
104     sym->etype->OCLASS = PRI   ;
105
106     var_dcl( yybss, PRI, sym, ";" );
107     yybss( "\t/* %s [%s(), static local] */\n", sym->name, funct_name );
108 }
109 /*-----*/
110
111 void    remove_symbols_from_table( sym )
112 symbol  *sym;
113 {
114     /* Remove all symbols in the list from the table. */
115
116     symbol *p;
117
118     for( p = sym; p ; p = p->next )
119         if( !p->duplicate )
120             delsym( Symbol_tab, p );
121         else
122         {
123             yyerror("INTERNAL, remove_symbol: duplicate sym. in cross-link\n");
124             exit( 1 );
125         }
126 }

```

Automatic variables:
`figure_local_offsets()`.

The `figure_local_offsets()` call on line 589 of Listing 6.61 handles automatic variables. (The subroutine is on line 45 of Listing 6.62.) It goes through the list, adjusting the rnames to hold a string which, when used in an operand, references the variable. Ultimately, the reference will look something like this: `WP(fp+6)`, but only the `fp+6` is created here—the `WP` and parentheses are added later by the expression-processing code.

`Offset`.

The current offset from the base of the automatic-variable region is remembered in `Offset`, which is incremented by the size of each variable as space for it is allocated.

The variables are packed as closely as possible into the local-variable space. If alignment permits, they are placed in adjacent bytes. As with structures, padding is inserted if necessary to guarantee alignment. `figure_local_offsets()` is called at the top of every block, but `Offset` is reset to zero only at the top of the subroutine, so the size of the local-variable region continues to grow over the life of the subroutine as automatic variables are allocated. The final size of the region is determined once the entire subroutine has been processed by calling `loc_var_space()` on line 33 of Listing 6.62. This value is used at the end of the subroutine-processing code to define one of the macros that's passed to the `link` instruction.

The `create_static_locals()` call on line 590 of Listing 6.61 handles `static` locals. The subroutine starts on line 84 of Listing 6.62. It goes through the list a second time, allocating space for the variable as if it were a `static` global variable. An arbitrary name is assigned instead of using the declared name, as was the case with true globals. This way, two subroutines can use the same name for a static variable without a conflict. All `static` locals are declared `private`, so conflicts with variables in other modules are not a problem. Symbols are removed from the table when the compiler leaves the scoping-level on lines 582 and 583 of Listing 6.61.

`loc_var_space()`.

Static-local variables:
`create_static_locals()`

6.6.6 Front-End/Back-End Considerations

Packing variables onto the stack makes sense in a one-pass compiler, but it's contraindicated if the real code is going to be generated by a back end. The front end has no way of knowing what alignment restrictions apply to the target machine or the actual sizes of the various types. The back end's life can be made easier by assuming that all types are the same size and that there are no alignment restrictions. As it is, the back end might have to undo some of our work. If an `int` is 32 bits, it must unpack the variables. Similarly, if the worst-case alignment restriction is two rather than four, it must get rid of the extra padding.

Don't pack variables if back end is used.

If the front end is ignoring size and alignment, some mechanism is needed to pass the symbol-table information to the back end. Currently, the compiler's just throwing that information away when it leaves the scoping level. A better approach passes the symbol table to the back end as part of the intermediate code. For example, you can introduce a new, local storage class to C-code and generate definitions for all local symbols at the top of a block along with the `static`-variable definitions. A matching `delete(name)` directive can be generated at the end of the block to tell the back end that the symbol had gone out of scope.

Passing symbol table information to back end.

There's no need to worrying about the size at the intermediate-code level if all variables are the same size. The compiler's been keeping track of the size so that we can use `W()`, `L()` and so forth to access variables, but it wouldn't have to do so if everything was the same size. Consequently, both global and local variables can just be called out by name in the intermediate code. You can dispense with all the size-related, C-code addressing modes and register-access directives and just use the names: `_p`, rather than `WP(&_p)` or `WP(fp-16)`, `r0` rather than `r0.pp`. The back end can compute the offsets for local variables and make any necessary adjustments to the generated code, replacing the symbolic names with stack-relative access directives as necessary. This one change dramatically simplifies both the design of C-code and the complexity of the front end.

No need to keep track of sizes.

All of the foregoing applies to structures as well as simple variables—the fields should all be the same size and not be packed into the structure. Better yet, the front end could make no attempt to determine the offset to the field from the base of the structure. Structure members could be passed to the back end like this:

`member type structure_name.field_name;`

and the fields could be called out by name in the intermediate code rather than generating explicit offsets to them, using something like `struct_name.member_name`.

6.7 The `gen()` Subroutine

Reasons to use `gen()`.

The `gen()` subroutine was used in the last section to print out the few C-code instructions in the subroutine prefix and suffix. `gen()` is a general-purpose code-generation interface for the parser—all C-code instructions are emitted using `gen()` calls rather than `yycode()` calls. It seems reasonable to look at it now, before using it further. I've concentrated the code emission into a single subroutine for several reasons:

- Clarity in the source code. Once you've added leading tabs, trailing newlines, field widths, and so forth, direct `yycode()` calls are pretty hard to read. Since `gen()` takes care of all the formatting for you, the subroutine calls are more understandable, and the code more maintainable as a consequence. `gen()` also makes the output code more readable because the code is formatted consistently.
- Fewer C-code syntax errors. There are slight variations in syntax in the C-code instruction set. Some instructions must be followed by semicolons, others by colons, and still others by no punctuation at all. Some instructions take parenthesized arguments, others do not. `gen()` takes care of all these details for you, so the odds of a syntax error showing up in the output are much smaller.
- Portability. Since all the output is concentrated in one place, it's much easier to make changes to the intermediate language. You need only change a single subroutine instead of several `yycode()` calls scattered all over the parser. Similarly, it's easy to emit binary output rather than ASCII output—just change `gen()` to emit binary directly.¹⁹
- Debugging. The compiler takes a command line switch that causes it to generate a run-time trace. Instead of emitting a single C-code directive, the compiler emits the C-code directive surrounded by statements that print the directive itself, the contents of all the virtual registers, and the top few stack elements. This way, you can watch the effect of every output instruction as it's executed. It is much easier to emit these extra run-time-trace statements when all output is concentrated in one place.

@ as first character in format string.

The first argument to `gen()` is a string that specifies the instruction to emit—usually the op code. The number and type of any additional arguments are controlled by the first one—legal first arguments are summarized in Table 6.16. If an argument is a character-pointer, the string is printed; if it's an `int`, the number is converted to a string and printed, and so on. In addition, if the first character of an arithmetic instruction is an `@`, the `@` is removed and a `*` is printed to the left of the destination string. This call:

```
gen( "@+=", "dst", "src" );
```

generates this code:

```
*dst += src;
```

Add comments to output code, `gen_comment()`.

The `gen()` subroutine is implemented in Listing 6.63 along with various support routines. `gen_comment()` (on line 81) puts comments in the output. It works like `printf()`, except that the output is printed to the right of the instruction emitted by the

19. This argument also holds for the code that creates declarations. I should have funneled all declarations through a single subroutine rather than using direct `yydata()` and `yybss()` calls. You may want to make that change to the earlier code as an exercise.

Table 6.16. The gen() Interface

First Argument	Second Argument	Third Argument	Output	Description
"%"	char *dst;	char *src;	dst %= src;	modulus
"&="	char *dst;	char *src;	dst &= src;	bitwise AND
"*="	char *dst;	char *src;	dst *= src;	multiply
"*=%s%d"	char *dst;	int src;	dst *= src;	multiply dst by constant
"+="	char *dst;	char *src;	dst += src;	add
"+=%s%d"	char *dst;	int src;	dst += src;	add constant to dst
"-="	char *dst;	char *src;	dst -= src;	subtract
"-=%s%d"	char *dst;	int src;	dst -= src;	subtract constant from dst
"/="	char *dst;	char *src;	dst /= src;	divide
"/=%s%d"	char *dst;	int src;	dst /= src;	divide dst by constant
"<=="	char *dst;	char *src;	dst <= src;	left shift dst by src bits
">=="	char *dst;	char *src;	dst >= src;	right shift dst by src bits
">L="	char *dst;	char *src;	lrs(dst,src);	logical right shift dst by src bits
"=-"	char *dst;	char *src;	dst =- src;	two's complement
"=~"	char *dst;	char *src;	dst =~ src;	one's complement
" ="	char *dst;	char *src;	dst = src;	bitwise OR
"^="	char *dst;	char *src;	dst ^= src;	bitwise XOR
"="	char *dst;	char *src;	dst = src;	assign
"=&"	char *dst;	char *src;	dst = &src;	load effective address
"=*%s%v"	char *dst;	value *src;	dst = *name; dst = name;	assign indirect. name is taken from src->name. If the name is of the form &name, then dst=name is output, otherwise dst=*name is output.
:	char *label;	(none)	label:	label
":%s%d"	char *alpha;	int num;	alphanum:	label label, but with the alphabetic and numeric components specified separately. gen(":%s%d", "P", 10) emits P10:.
"BIT"	char *opl;	char *bit;	BIT(opl,bit)	test bit
"EQ"	char *opl;	char *op2;	EQ(opl,op2)	equality
"EQ%s%d"	char *opl;	int op2;	EQ(opl,op2)	equal to constant
"GE"	char *opl;	char *op2;	EQ(opl,op2)	greater than or equal
"GT"	char *opl;	char *op2;	EQ(opl,op2)	greater than
"LE"	char *opl;	char *op2;	EQ(opl,op2)	less than or equal
"LT"	char *opl;	char *op2;	EQ(opl,op2)	less than
"NE"	char *opl;	char *op2;	EQ(opl,op2)	not equal
"U_EQ"	char *opl;	char *op2;	EQ(opl,op2)	greater than or equal, unsigned
"U_GT"	char *opl;	char *op2;	EQ(opl,op2)	greater than, unsigned
"U_LE"	char *opl;	char *op2;	EQ(opl,op2)	less than or equal, unsigned
"U_LT"	char *opl;	char *op2;	EQ(opl,op2)	less than, unsigned
"PROC"	char *name;	char *cls;	PROC(name,cls)	start procedure
"ENDP"	char *name;	(none)	ENDP(name)	end procedure
"call"	char *label;	(none)	call(label);	call procedure
"ext_high"	char *dst;	(none)	ext_high(dst);	sign extend
"ext_low"	char *dst;	(none)	ext_low(dst);	sign extend
"ext_word"	char *dst;	(none)	ext_word(dst);	sign extend
"goto"	char *label;	(none)	goto label;	unconditional jump
"goto%s%d"	char *alpha;	int num;	goto alphanum;	unconditional jump, but the alphabetic and numeric components of the target label are specified separately. gen("goto%s%d", "P", 10) emits goto P10;.
"link"	char *loc;	char *tmp;	link(loc+tmp);	link
"pop"	char *dst;	char *type;	dst = pop(type);	pop
"push"	char *src;	(none)	push(src);	push
"ret"	(none)	(none)	ret();	return
"unlink"	(none)	(none)	ret();	unlink

next gen() call. gen_comment() stores the comment text in Comment_buf (declared on line 77) so that it can be printed by a subsequent gen() call.

The enable_trace() and disable_trace() subroutines on lines 100 and 101 enable and disable the generation of run-time trace statements.

Run-time trace,
enable_trace(),
disable_trace().

Strength reduction.

Instruction output and
formatting:

`print_instruction()`

Run-time trace instruc-
tions: `_P()`, `_T()`.

`gen()` itself starts on line 113. It uses the lookup table on lines 21 to 74 to translate the first argument into one of the tokens defined on lines 13 to 19. The table lookup is done by the `bsearch()` call on line 132. Thereafter, the token determines the number and types of the arguments, which are pulled off the stack on lines 138 to 149. The ANSI variable-argument mechanism described in Appendix A is used. `gen()` does not emit anything itself—it assembles a string which is passed to a lower-level output routine. The switch starting on line 156 takes care of most of the formatting, using `sprintf()` calls to initialize the string. Note that a simple optimization, called *strength reduction* is done on lines 201 to 229 in the case of multiplication or division by a constant. If the constant is an even power of two, a shift is emitted instead of a multiply or divide directive. A multiply or divide by 1 generates no code at all. The comment, if any, is added to the right of the output string on line 236.

The actual output is done in `print_instruction()` on line 247 of Listing 6.63. This subroutine takes care of all the formatting details: labels are not indented, the statement following a test is indented by twice the normal amount, and so forth. `print_instruction()` also emits the run-time trace directives. These directives are written directly to `yycodeout` [with `fprintf()` calls rather than `yycode()` calls] so that they won't show up in the IDE output window. The trace is done using two macros: `_P()` prints the instruction, and `_T()` dumps the stack and registers. Definitions for these macros are written to the output file the first time that `print_instruction()` is called with tracing enabled (on Line 262 of Listing 6.63). The output definitions look like this:

```
#define _P(s) printf( s )
#define _T() pm(),printf("-----\\n")
```

`_P()` just prints its argument, `_T()` prints the stack using a `pm()` call—`pm()` is declared in

Most statements are handled as follows:

```
a = b;           _P( "a = b;" )  
                _T()
```

The trace directives are printed at the right of the page so that the instructions themselves will still be readable. The instruction is printed first, then executed, and then the stack and registers are printed. Exceptions to this order of events are as follows:

label:	<code>_P("label:")</code>
PROC(...)	<code>_P("PROC(...)"); _T();</code>
	<code>_P("ret(...)")</code>
ret(...);	
	<code>_P("ENDP(...)")</code>
ENDP(...)	

The trace for a logical test is tricky because both the test and the following instruction must be treated as a unit. They are handled as follows:

```
NE(a,b)           _P( "NE(a,b)" )  
                 {  
instruction;     _P( "goto x" );  
                 _T(); }
```

Listing 6.63. gen.c—C-code Generation

```

1 #include <stdio.h>
2 #include <stdarg.h>
3 #include <tools/debug.h>
4 #include <tools/hash.h>
5 #include <tools/l.h>
6 #include <tools/compiler.h>
7 #include "symtab.h"
8 #include "value.h"
9 #include "proto.h"
10
11 PRIVATE int Trace = 0;           /* Generate run-time trace if true. */
12
13 typedef enum request
14 {
15     t_assign_addr, t_assign_ind, t_call, t_endp, t_ext, t_goto, t_goto_int,
16     t_label, t_label_int, t_link, t_logical, t_logical_int, t_lrs, t_math,
17     t_math_int, t_pop, t_proc, t_push, t_ret, t_unlink
18 }
19 } request;
20
21 struct ltab
22 {
23     char      *lexeme;
24     request  token;
25 }
26 Ltab[] =
27 {
28     {"%=",          t_math          },
29     {"&=",          t_math          },
30     {"*=",          t_math          },
31     {"*=%s%d",      t_math_int     }, /* Multiply var by constant. */
32     {"+=" ,         t_math          },
33     {"+=%s%d",      t_math_int     },
34     {"-=",          t_math          },
35     {"-=%s%d",      t_math_int     },
36     {"/=",          t_math          },
37     {"/==%s%d",      t_math_int     },
38     {":",            t_label         },
39     {":%s%d",       t_label_int    },
40     {"<=>=",       t_math          },
41     {"==",          t_math          },
42     {"=&",         t_assign_addr   }, /* Get effective address. */
43     {"=*%s%v",       t_assign_ind   }, /* Assign indirect. */
44     {"=-",          t_math          },
45     {">>=",       t_math          },
46     {">L=",        t_lrs           },
47     {"BIT",          t_logical       },
48     {"ENDP",         t_endp          },
49     {"EQ",           t_logical       },
50     {"EQ%s%d",      t_logical_int  },
51     {"GE",           t_logical       },
52     {"GT",           t_logical       },
53     {"LE",           t_logical       },
54     {"LT",           t_logical       },
55     {"NE",           t_logical       },
56     {"PROC",         t_proc          },
57     {"U_GE",         t_logical       },
58     {"U_GT",         t_logical       },
59     {"U_LE",         t_logical       },

```



Listing 6.63. continued...

```

60      {"U_LT",           t_logical        },
61      {"^=",             t_math           },
62      {"call",            t_call           },
63      {"ext_high",        t_ext            },
64      {"ext_low",          t_ext            },
65      {"ext_word",         t_ext            },
66      {"goto",             t_goto           },
67      {"goto%s%d",        t_goto_int       },
68      {"link",             t_link           },
69      {"pop",              t_pop            },
70      {"push",             t_push           },
71      {"ret",              t_ret            },
72      {"unlink",           t_unlink          },
73      {"|=",               t_math           }
74 };
75
76 #define NREQ ( sizeof(Ltab)/sizeof(*Ltab) ) /* Table size (in elements). */
77 char Comment_buf[132];                      /* Remember comment text here. */
78
79 /*-----*/
80
81 PUBLIC void gen_comment( format, ... )
82 char *format;
83 {
84     /* Works like printf(), but the string is appended as a comment to the end
85      * of the command generated by the next gen() call. There's no array-size
86      * checking---be careful. The maximum generated string length is 132
87      * characters. Overwrite any comments already in the buffer.
88     */
89
90     va_list args;
91     va_start ( args, format );
92     vsprintf ( Comment_buf, format, args );
93     va_end   ( args );
94 }
95
96 /*-----*/
97     /* Enable/disable the generation of run-time trace output.
98 */
99
100 PUBLIC enable_trace() { Trace = 1; } /* Must call before parsing starts. */
101 PUBLIC disable_trace() { Trace = 0; }
102
103 /*-----*/
104
105 PRIVATE int cmp( a, b ) /* Compare two lexeme fields of an ltab. */
106 struct ltab *a, *b;
107 {
108     return strcmp( a->lexeme, b->lexeme );
109 }
110
111 /*-----*/
112
113 PUBLIC gen( op ,... )                  /* emit code */
114 char *op;
115 {
116     char      *dst_str, *src_str, b[80];
117     int       src_int, dst_int;
118     value    *src_val;
119     struct ltab *p, dummy;

```



Listing 6.63. continued...

```

120     request      tok;
121     va_list      args;
122     char          *prefix = " ";
123     int           amt;
124
125     if( *op == '@' )
126     {
127         ++op;
128         prefix = "**";
129     }
130
131     dummy.lexeme = op;
132     if ( ! (p = (struct ltab *) bsearch(&dummy, Ltab, NREQ, sizeof(*Ltab), cmp)))
133     {
134         yyerror("INTERNAL gen: bad request <%s>, no code emitted.\n", op );
135         return;
136     }
137
138     va_start( args, op );                                /* Get the arguments. */
139     dst_str = va_arg( args, char* );
140     switch( tok = p->token )
141     {
142         case t_math_int:
143         case t_logical_int:
144         case t_goto_int:
145         case t_label_int:   src_int = va_arg( args, int ); break;
146         case t_assign_int: src_val = va_arg( args, value* ); break;
147         default:          src_str = va_arg( args, char* ); break;
148     }
149     va_end( args );
150
151     /* The following code just assembles the output string. It is printed with
152      * the print_instruction() call under the switch, which also takes care of
153      * inserting trace directives, inserting the proper indent, etc.
154      */
155
156     switch( tok )
157     {
158         case t_call:        sprintf(b, " call(%s);",    dst_str ); break;
159         case t_endp:       sprintf(b, " ENDP(%s)",    dst_str, src_str ); break;
160         case t_text:        sprintf(b, " %s(%s);",    op,    dst_str ); break;
161         case t_goto:       sprintf(b, " goto %s;",    dst_str ); break;
162         case t_goto_int:   sprintf(b, " goto %s%d;", dst_str, src_int ); break;
163         case t_label:      sprintf(b, "%s:",        dst_str ); break;
164
165         case t_label_int:  sprintf(b, "%s%d:",      dst_str, src_int );
166             tok = t_label;
167             break;
168
169         case t_logical:    sprintf(b, " %s(%s,%s)", op,    dst_str, src_str );
170             break;
171
172         case t_logical_int: sprintf(b, " %2.2s(%s,%d)", op, dst_str, src_int );
173             tok = t_logical;
174             break;
175
176         case t_link:       sprintf(b, " link(%s+%s);", dst_str, src_str ); break;
177         case t_pop:        sprintf(b, " %-12s = pop(%s);",dst_str, src_str ); break;

```



Listing 6.63. continued...

```

178     case t_proc:      sprintf(b," PROC(%s,%s)",      dst_str, src_str ); break;
179     case t_push:      sprintf(b," push(%s);",        dst_str           ); break;
180     case t_ret:       sprintf(b," ret();",          dst_str           ); break;
181     case t_unlink:    sprintf(b," unlink();" ,       dst_str           ); break;
182
183     case t_lrs:       sprintf(b,"%slrs(%s,%s);",   prefix, dst_str, src_str );
184             break;
185
186     case t_assign_addr: sprintf(b,"%s%-12s = &%s;", prefix, dst_str, src_str );
187             break;
188     case t_assign_ind:
189         if( src_val->name[0]=='&' )
190             sprintf(b,"%s%-12s = %s;",prefix, dst_str, src_val->name+1);
191         else
192             sprintf(b,"%s%-12s = *%s;",prefix, dst_str, src_val->name);
193
194             break;
195
196     case t_math:      sprintf(b,"%s%-12s %s %s;", prefix, dst_str, op, src_str );
197             break;
198     case t_math_int:
199         if( *op != '*' && *op != '/' )
200             sprintf(b,"%s%-12s %2.2s %d;", prefix, dst_str, op, src_int);
201         else
202         {
203             switch( src_int )
204             {
205                 case 1 :   amt = 0;      break;
206                 case 2 :   amt = 1;      break;
207                 case 4 :   amt = 2;      break;
208                 case 8 :   amt = 3;      break;
209                 case 16:  amt = 4;      break;
210                 case 32:  amt = 5;      break;
211                 case 64:  amt = 6;      break;
212                 case 128: amt = 7;      break;
213                 case 256: amt = 8;      break;
214                 case 512: amt = 9;      break;
215                 case 1024: amt = 10;    break;
216                 case 2048: amt = 11;    break;
217                 case 4096: amt = 12;    break;
218                 default: amt = -1;    break;
219             }
220
221             if( !amt )
222                 sprintf(b, "/* %s%-12s %s 1; */", prefix, dst_str, op );
223
224             else if( amt < 0 )
225                 sprintf(b, "%s%-12s %s %d;", prefix, dst_str, op, src_int);
226             else
227                 sprintf(b, "%s%-12s %s %d;", prefix, dst_str,
228                               (*op == '*') ? "<==" : ">==" , amt);
229             }
230             break;
231     default:
232         yyerror("INTERNAL, gen: bad token %s, no code emitted.\n", op );
233         break;
234     }
235

```



Listing 6.63. continued...

```

236     if( *Comment_buf )           /* Add optional comment at end of line. */
237     {
238         concat( sizeof(b), b, b, (tok == t_label ? "\t\t\t\t" : "\t"),
239                  "/* ", Comment_buf, " */", NULL );
240         *Comment_buf = '\0';
241     }
242     print_instruction( b, tok );      /* Output the instruction. */
243 }
244
245 /*-----*/
246
247 PRIVATE void print_instruction( b, t )
248 char    *b;          /* Buffer containing the instruction.      */
249 request t;          /* Token representing instruction.      */
250 {
251     /* Print the instruction and, if trace is enabled (Trace is true), print
252     * code to generate a run-time trace.
253     */
254
255     extern FILE  *yycodeout, *yybssout;
256     static int   printed_defs = 0;
257     static int   last_stmt_was_test;
258
259     if( Trace && !printed_defs )
260     {
261         printed_defs = 1;
262         fprintf( yybssout,
263                 "#define _P(s) printf( s )\n"
264                 "#define _T() pm(),printf(\"-----\\n\")"
265                 "\n\n" );
266     }
267
268     if( !Trace )                      /* just print the instruction */
269     {
270         yycode("%s%s%s\n", (t==t_label || t==t_endp || t==t_proc) ? "" : "\t",
271                 (last_stmt_was_test) ? " " : "", b);
272
273         last_stmt_was_test = (t==t_logical);
274     }
275     else if( t == t_logical )
276     {
277         fprintf( yycodeout, "\t\t\t\t\t" "_P(\"%s\\n\");\n", b);
278         yycode("\t%s\t\t{\n", b);                                /* */ */
279         last_stmt_was_test = 1;
280     }
281     else
282     {
283         switch( t )
284         {
285             case t_label: yycode("%s", b);
286                         fprintf( yycodeout, "\t\t\t\t\t" "_P(\"%s\\n\");", b);
287                         break;
288
289             case t_proc:  yycode("%s", b);
290                         fprintf( yycodeout, "\t\t\t\t" "_P(\"%s\\n\");", b);
291                         fprintf( yycodeout, "\t\t\t\t" "_T();", );
292                         break;
293
294

```

Listing 6.63. continued...

6.8 Expressions

This section looks at how expressions are processed. Temporary-variable management is discussed as are lvalues and rvalues. The code-generation actions that handle expression processing are covered as well. We've looked at expression parsing sufficiently that there's no point in including extensive sample parses of every possible production in the current section—I've included a few sample parses to explain the harder-to-understand code-generation issues, but I expect you to be able to do the simpler cases yourself. Just remember that order of precedence and evaluation controls the order in which the productions that implement particular operators are executed. Productions are reduced in the same order as the expression is evaluated.

The foregoing notwithstanding, if you have the distribution disk, you may want to run simple expressions through the compiler (*c.exe*) as you read this and the following sections. As you watch the parse, pay particular attention to the order in which reductions occur and the way that attributes are passed around as the parse progresses.

6.8.1 Temporary-Variable Allocation

All expressions in C are evaluated one operator at a time, with precedence and associativity determining the order of evaluation as much as is possible. It is conceptually convenient to look at every operation as creating a temporary variable that somehow references the result of that operation—our compiler does things somewhat more efficiently, but it's best to think in terms of the stupidest possible code. This temporary, which represents the evaluated subexpression, is then used as an operand at the next expression-evaluation stage.

Defer temporary-variable management to back end.

Our first task is to provide a mechanism for creating and deleting temporaries. One common approach is to defer the temporary-variable management to the back end. The compiler itself references the temporaries as if they existed somewhere as global variables, and the back end takes over the allocation details. The temporary variable's type can be encoded in the name, using the same syntax that would be used to access the temporary from a C-code statement: `W(t0)` is a word, `L(t1)` is an lword, `WP(t2)` is a word pointer, and so on. The advantage of this approach is that the back end is in a

much better position than the compiler itself to understand the limitations and strengths of the target machine, and armed with this knowledge it can use registers effectively. I am not taking this approach here because it's pedagogically useful to look at a worst-case situation—where the compiler itself has to manage temporaries.

Temporary variables can be put in one of three places: in registers, in static memory, and on the stack. The obvious advantage of using registers is that they can be accessed quickly. The registers can be allocated using a stack of register names—essentially the method that is used in the examples in previous chapters. If there aren't enough registers in the machine, you can use a few run-time variables (called *pseudo registers*) as replacements, declaring them at the top of the output file and using them once the registers are exhausted. You could use two stacks for this purpose, one of register names and another of variable names, using the variable names only when there are no more registers available. Some sort of priority queue could also be used for allocation.

Registers as temporaries, pseudo registers.

One real advantage to deferring temporary-variable allocation to the back end is that the register-versus-static-memory problem can be resolved in an efficient way. Many optimizers construct a syntax tree for the expression being processed, and analysis of this tree can be used to allocate temporaries efficiently (so that the registers are used more often than the static memory). This sort of optimization must be done by a postprocessor or postprocessing stage in the parser, however—the parser must create a physical syntax or parse tree that a second pass can analyze. Though it's easy for a simple one-pass compiler to use registers, it's difficult for such a compiler to use the registers effectively.

Problems with function calls.

Another problem is function calls, which can be imbedded in the middle of expressions. Any registers or pseudo registers that are in use as temporaries must be pushed before calling the function and popped after the return. Alternately, code at the top of the called function could push only those registers that are used as temporaries in the function itself—there's a lot of pushing and popping in either case. This save-and-restore process adds a certain amount of overhead, at both compile time and run time.

Temporaries on stack.

The solution to the temporary-variable problem that's used here is a compromise between speed and efficiency. A region of the stack frame is used for temporaries.²⁰ This way, they don't have to be pushed because they're already on the stack.

Because the maximum size of the temporary-variable region needed by a subroutine varies (it is controlled by the worst-case expression in the subroutine), the size of the temporary-variable space changes from subroutine to subroutine. This problem is solved with the second macro that's used in the `l1nk` instruction in the subroutine prefix. (`L1` in the example in the previous section.) The macro is defined to the size of the temporary-variable region once the entire subroutine has been processed.

Dynamic temporary-variable creation.

This approach—allocating a single, worst-case sized temporary-variable region—is generally better than a dynamic approach where the temporary-variable space is gradually expanded at run time by subtracting constants from the stack pointer as variables are needed. The stack is shrunk with matching additions when the variable is no longer needed. This last approach can be more efficient of stack space, but it is both more difficult to do at compile time and is inefficient at run time because several subtractions and additions are needed rather than a single `l1nk` instruction. Since most languages use very few, relatively small, temporaries as they evaluate expressions, this second

20. In many machines, such as the Intel 8086 family, a stack-relative memory access is actually more efficient than a direct-mode memory access. It takes fewer clock cycles. Since none of the 8086-family machines have any general purpose registers to speak of, putting the temporaries on the stack is actually one of the most efficient solutions to the problem.

method is usually more trouble than its worth. Nonetheless, you may want to consider a dynamic approach if the source language has very large data types. For example, a special-purpose language that supported a matrix basic type and a matrix-multiplication operator would need a tremendous amount of temporary-variable space to process the following expression:

```
matrix a[100][100], b[100][100], c[100][100];  
a = a * b * c;
```

It would be worthwhile not to waste this stack space except when the expression was actually being evaluated, so a dynamic approach makes sense here. If the stack just isn't large enough for temporaries such as the foregoing, the compiler could generate run-time calls to `malloc()` as part of the subroutine prefix, storing a pointer to the memory returned from `malloc()` on the stack. A run-time `free()` call would have to be generated as part of the subroutine suffix to delete the memory. Calling `malloc()` is a very-high-overhead operation, however, so it's best to use the stack if you can.

Temporary variable's type.

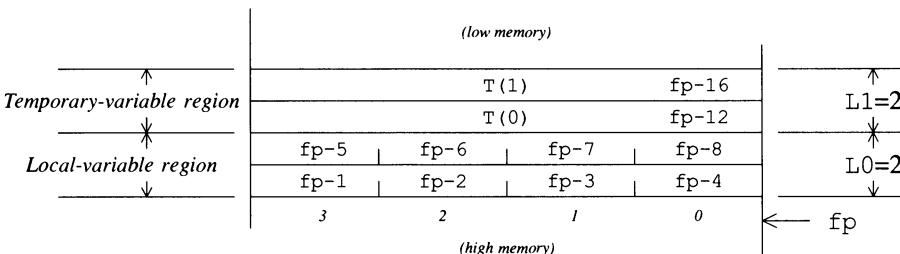
The next problem is the temporary variable's type, which varies with the operand types. The type of the larger operand is used for the temporary. The easiest solution (and the one used here) is to use the worst-case size for all temporaries, and align them all on a worst-case boundary. All temporaries take up one 1word-sized stack element, which is guaranteed to be aligned properly because it's a stack word. Again, if the worst-case type was very large (if doubles were supported, for example), it would be worthwhile to go to the trouble to pack the temporaries into the stack region in a manner analogous to the local variables.

Figure 6.18 shows a stack frame with eight bytes of local-variable space and enough temporary space for two variables. The macro definitions at the top of the figure are generated at the same time as the subroutine prefix. The `L0` macro evaluates to the size, in stack elements, of the local-variable region. (This is the same `L0` that's passed to the link instruction.) The `T()` macro simplifies temporary-variable access in instructions. It is defined at the top of the file as an empty macro, and is then `#undefined` and redefined to reference the current local-variable-space size at the top of every subroutine. The initial dummy definition is required because some compilers won't let you undefine a nonexistent macro. The following examples demonstrate how the macro is used:

<code>W(T(0))</code>	accesses the word in the bottom two bytes of the bottom temporary.
<code>L(T(1))</code>	accesses the lword that takes up the entire top temporary.
<code>*WP(T(0))</code>	accesses the word whose address is in the top temporary.

Figure 6.18. Storing Temporary Variables in the Stack Frame

```
#define T(n)    (fp-L0-(4*n))
#define L0     2           /* size of local-variable space      */
#define L1     2           /* size of temporary-variable space */
```



Though temporary variables are not packed into the stack frame, the temporary-variable space is used intelligently by the compiler—the same region of the stack is recycled as needed for temporaries. The strategy, implemented in Listing 6.64, is straightforward. (It's important not to confuse compile and run time, in the following discussion. It's easy to do. The temporary-variable allocator is part of the compiler, but it is keeping track of a condition of the stack at run time. All of the temporary-variable allocation and deallocation is done at compile time. A region of the stack is just used as a temporary at run time—the compiler determined that that region was available at compile time.)

Recycling temporaries.

Listing 6.64. *temp.c*—Temporary-Variable Management

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/debug.h>
4 #include <tools/hash.h>
5 #include <tools/l.h>
6 #include <tools/compiler.h>
7 #include <tools/c-code.h>
8 #include "symtab.h"
9 #include "value.h"
10 #include "proto.h"
11
12 /* Subroutines in this file take care of temporary-variable management. */
13
14 #define REGION_MAX 128 /* Maximum number of stack elements that can be */
15 /* used for temporaries. */
16 #define MARK -1 /* Marks cells that are in use but are not the */
17 /* first cell of the region. */
18 typedef int CELL; /* In-use (Region) map is an array of these. */
19
20 PRIVATE CELL Region[REGION_MAX];
21 PRIVATE CELL *High_water_mark = Region;
22 /*-----*/
23
24 PUBLIC int tmp_alloc( size )
25 int size; /* desired number of bytes */
26 {
27     /* Allocate a portion of the temporary-variable region of the required size,
28      * expanding the tmp-region size if necessary. Return the offset in bytes
29      * from the start of the rvalue region to the first cell of the temporary.
30      * 0 is returned if no space is available, and an error message is also
31      * printed in this situation. This way the code-generation can go on as if
32      * space had been found without having to worry about testing for errors.
33      * (Bad code is generated, but so what?) */
34 }
35
36 CELL *start, *p;
37 int i;
38
39 /* size = the number of stack cells required to hold "size" bytes. */
40
41 size = ((size + SWIDTH) / SWIDTH) - (size % SWIDTH == 0);
42
43 if( !size )
44     yyerror("INTERNAL, tmp_alloc: zero-length region requested\n");
45
46

```



Listing 6.64. continued...

```

47     /* Now look for a large-enough hole in the already-allocated cells. */
48
49     for( start = Region; start < High_water_mark ; )
50     {
51         for( i = size, p = start; --i >= 0 && !*p; ++p )
52             ;
53
54         if( i >= 0 )                                /* Cell not found. */
55             start = p + 1;
56         else                                         /* Found an area big enough. */
57             break;
58     }
59
60     if( start < High_water_mark )                /* Found a hole. */
61         p = start;
62     else
63     {
64         if( (High_water_mark + size) > (Region + REGION_MAX) )    /* No room. */
65         {
66             yyerror("Expression too complex, break into smaller pieces\n");
67             return 0;
68         }
69         p = High_water_mark;
70         High_water_mark += size;
71     }
72
73     for( *p = size; --size > 0; *++p = MARK) /* 1st cell=size. Others=MARK */
74     ;
75
76     return( start - Region );                  /* Return offset to start of region */
77                                         /* converted to bytes. */
78 }
79
80 /*-----*/
81
82 PUBLIC void tmp_free( offset )
83 int offset;                               /* Release a temporary var.; offset should */
84 {                                         /* have been returned from tmp_alloc(). */
85     CELL *p = Region + offset;
86     int size;
87
88     if( p < Region || p > High_water_mark || !*p || *p == MARK )
89         yyerror( "INTERNAL, tmp_free: Bad offset (%d)\n", offset );
90     else
91         for( size = *p; --size >= 0; *p++ = 0 )
92             ;
93 }
94
95 /*-----*/
96
97 PUBLIC void tmp_reset()
98 {
99     /* Reset everything back to the virgin state, including the high-water mark.
100      * This routine should be called just before a subroutine body is processed,
101      * when the prefix is output. See also: tmp_freeall().
102      */
103     tmp_freeall();
104     High_water_mark = Region ;
105 }

```



Listing 6.64. continued...

```

107  /*-----*/
108
109 PUBLIC void tmp_freeall()
110 {
111     /* Free all temporaries currently in use (without modifying the high-water
112      * mark). This subroutine should be called after processing arithmetic
113      * statements to clean up any temporaries still kicking around (there is
114      * usually at least one).
115      */
116     memset( Region, 0, sizeof(Region) );
117 }
118
119 /*-----*/
120
121 PUBLIC int tmp_var_space()
122 {
123     /* Return the total cumulative size of the temporary-variable region in
124      * stack elements, not bytes. This number can be used as an argument to the
125      * link instruction.
126      */
127     return High_water_mark - Region;
128 }
```

The compiler keeps an internal array called `Region`—declared on line 20 of Listing 6.64—that tells it which stack elements in the temporary-variable space are in use at run time. The index of the cell in the array corresponds to the offset of the equivalent stack element. For example, `Region[N]` corresponds to the temporary variable at offset N (in stack elements) from the base of the temporary-variable region (not from the frame pointer). The array elements are zero if the corresponding stack element is available. For example, if `Region[4]` is zero, then the stack element at offset 4 from the base of the temporary space is available. The size of the array is controlled by `REGION_MAX`, declared on line 14 of Listing 6.64. This size limits the amount of temporary-variable space available to a subroutine. This size limits an expression’s complexity because it limits the number of temporaries that can be used to evaluate the expression. The compiler prints an error message if the expression gets too complicated (if there aren’t enough temporary variables to evaluate all the subexpressions).

Compile-time map of
run-time stack: `Re-
gion[]`.

The temporary-variable allocator, `tmp_alloc()`, starts on line 24 of Listing 6.64. It is passed the number of bytes required for the temporary, figures out how many stack elements are required for that temporary, and then allocates that many bytes from the temporary-variable region.²¹ The allocation logic on lines 49 to 58 of Listing 6.64 searches the array from bottom to top looking for a free stack element. It allocates the first available stack element for the temporary. The `High_water_mark`, used here and declared on line 21 of Listing 6.64, keeps track of the size of the temporary-variable space. If the first available element is at a higher index than any previously allocated element, then `High_water_mark` is modified to point at the new element.

Allocate temporary:
`tmp_alloc()`.

`High_water_mark`.

21. The allocator I’m presenting here is actually more complex than necessary, because all temporary variables take up exactly one stack element in the current compiler. It seemed reasonable to demonstrate how to handle the more complex situation, where some temporaries can take up more than one stack element, however.

Marking a stack cell as ‘in use’.

The cell is marked as “in use” on line 73. The `Region` element corresponding to the first cell of the allocated space is set to the number of stack elements that are being allocated. If more than one stack element is required for the temporary, adjacent cells that are part of the temporary are filled with a place marker. Other subroutines in Listing 6.64 de-allocate a temporary variable by resetting the equivalent `Region` elements to zero, de-allocate all temporary variables currently in use, and provide access to the high-water mark. You should take a moment and review them now.

6.8.2 Lvalues and Rvalues

Expression evaluation, and the way that temporaries are used while doing the evaluation, is actually a more complex proposition than you would think, based on the examples in earlier chapters. Hitherto, we’ve used temporaries for one purpose only, to hold the value of a subexpression. We could do this because none of the grammars that we’ve seen are capable of modifying variables, only of using them. As soon as modification becomes a possibility, things start to get complicated—especially in a language like C that provides several ways to modify something (`=`, `++`, `--`, and so forth).

Rvalues.

So far, the parsers that we’ve looked at always treated identifier references in the same way: They copied the contents of the referenced variable into a temporary and passed the name of that temporary back up as an attribute. The following grammar is a case in point:

```
e : e PLUS e { yycode("%s += %s\n", $1, $3); free_name($3); }
| ID          { yycode("%s = _%s\n", $$ = new_name(), yytext); }
;
```

The value stack is a stack of character pointers. The attribute is the name of the temporary that holds the partially-evaluated subexpression. Input like `x+y` generates the following code:

```
t0 = x;
t1 = y;
t1 += t0;
```

The kind of temporary that’s used here—which holds something’s value—is called an *rvalue*. In the current grammar, the `y` and `z` in `x=y+z` are said to *generate rvalues*—they create temporaries that hold a variable’s value. (It’s an *rvalue* because it can go to the right of the equals sign.) The `+` operator also generates an *rvalue* because the result of the operation is a run-time temporary holding the sum—the value of the subexpression.

Lvalues.

What happens when you introduce an assignment operator to the grammar? Our first, unsuccessful attempt at a solution treats the assignment operator just like the addition operator:

```
e : e PLUS e { yycode("%s += %s\n", $1, $3); free_name($3); }
| e EQUAL e { yycode("%s = %s\n", $1, $3); free_name($3); }
| ID          { yycode("%s = _%s\n", $$ = new_name(), yytext); }
;
```

The input `x=y+z` yields the following output:

```
t0 = x;
t1 = y;
t2 = z;
t1 += t2;
t0 = t1;
```

which doesn’t do anything useful.

In order to do assignment correctly, you need to know the address of `x` (the memory location to modify), not its contents. To get this information you need to introduce a second type of temporary—one that holds an address. This kind of temporary is called an *lvalue* because it can be generated by subexpressions to the *left* of an equal sign. The grammar and code in Listing 6.65 supports both kinds of temporaries. They are distinguished from one another by the first character of the name. Rvalues start with `R`, lvalues with `L`. As was the case in earlier chapters, a simple stack strategy is used to allocate and free names. Error checking has been omitted to simplify things.

Listing 6.65. Expression Processing with Lvalues and Rvalues

```

1  %{
2      #define YYSTYPE char*
3  %
4  %%
5  s      : e ;
6  e      : e PLUS e
7      {
8          char *target;
9
10         /* Convert $1 to an rvalue if necessary so that it can be used
11            * as the synthesized attribute.
12            */
13
14         if( *$1 != 'R' )
15         {
16             yycode( "%s = %s\n", target = new_rvalue(), rvalue($1) );
17             free_value($1);
18             $1 = target;
19         }
20         $$ = $1;
21         yycode("%s += %s\n", target, rvalue( $3 ) );
22         free_value( $3 );
23     }
24     | e EQUAL e
25     {
26         char *target;
27
28         if( *$1 != 'L' )
29             yyerror("Lvalue required\n");
30         else
31         {
32             if( *$3 != 'R' )           /* Convert $3 to an rvalue. */
33             {
34                 yycode( "%s = %s\n", target = new_rvalue(), rvalue($3) );
35                 free_value($3);
36                 $3 = target;
37             }
38             yycode( "*%s = %s", $1, $3 );
39             free_value( $1 );
40             $$ = $3;
41         }
42     }
43
44     | ID
45     {
46         yycode("%s = &_s\n", $$ = new_lvalue(), yytext );
47     }
48 ;
49 %%

```



Listing 6.65. continued...

```

50  char *Lvalues[4] = { "L0", "L1", "L2", "L3" };
51  char *Lp        = Lvalues;
52  char *Rvalues[4] = { "R0", "R1", "R2", "R3" };
53  char *Rp        = Rvalues;
54
55  char *new_lvalue() { return *Lp++; } /* Allocate a new lvalue. */
56  char *new_rvalue() { return *Rp++; } /* Allocate a new rvalue. */
57
58  void    free_value( val )           /* Free an lvalue or an rvalue. */
59  char    *val;
60 {
61     if( *val == 'L' ) --Lp = val;
62     else                --Rp = val;
63 }
64
65  char *rvalue( value )           /* If the argument is an lvalue, return a string */
66  char *value;                  /* that can be used in an expression to reference */
67 {                                /* the lvalue's contents. (Add a * to the name.) */
68     static char buf[16];
69
70     if( *lvalue == 'R' )          /* it's already an rvalue */
71         return value;
72     sprintf( buf, "*%s", lvalue );
73     return buf;
74 }
```

Using the new grammar, the input `x=y+z` yields the following output:

```

L0    = &_x;
L1    = &_y;
L2    = &_z;
R0    = *L1;
R0    += *L2;
*L0   = R0;
```

which is awkward, but it works. The lvalues are converted to rvalues by preceding their names with a star. Note that the attribute synthesized in the assignment is the rvalue, not the lvalue. The assignment operator *requires* an lvalue to the left of the equals sign, but it *generates* an rvalue. That's why code like the following is illegal in C:

`(x=y) = z;`

even though

`x = y = z;`

is okay. `(x=y)` generates an rvalue—a temporary that holds the value of `y`—and explicit assignment to an rvalue is meaningless. In `x=y=z`, the operand to the right of the equal sign is always an rvalue and the operand to the left is always an lvalue. (Run it through the parser by hand if you don't believe me.)

Logical versus physical
lvalues.

The code that the compiler just generated is pretty sorry. Fortunately, the compiler can do a simple optimization to fix things. Instead of passing around the name of a temporary variable as an attribute, it can pass a string which, when used as an operand in an instruction, evaluates correctly. This attribute is sometimes a temporary-variable name, but it can also take on other values. That is, instead of generating

`L0 = &_x;`

in the earlier example, and then passing around the string "`L0`" as an attribute, the

compiler can dispense with the assignment entirely and pass around "&_x" as an attribute. So you now have two types of lvalues: *physical* lvalues—temporary variables that hold the address of something, and *logical* lvalues—expressions that evaluate to the address of something when used as an operand. The compiler can distinguish between the two by looking at the first character in the name; logical lvalues start with an ampersand.

Physical and logical
values.

Listing 6.66 shows the parser, modified to use logical lvalues. Changes are marked with a ★. Very few changes have to be made. First, the code for assignment must be modified on line 32 of Listing 6.66 to handle logical lvalues. The conditional just discards the leading ampersand if it's present, because the two operators in the expression `* & _x` just cancel one another, yielding `_x`. Next, the `ID` action on lines 39 to 41 of Listing 6.66 must be changed to generate logical, rather than physical, lvalues. The name is assembled into a string, and a copy of the string is passed as an attribute. (The copy is freed on line 59.) Finally, the `rvalue()` subroutine is modified on lines 70 and 71 of Listing 6.66 to recognize both kinds of lvalues.

Listing 6.66. Expression Processing with Logical Lvalues

```

1 %{
2 #define YYSTYPE char*
3 %
4 %%
5 s      : e ;
6 e      : e PLUS e
7 {
8         char *target;
9
10        if( *$1 != 'R' )
11        {
12            yycode( "%s = %s\n", target = new_rvalue(), rvalue($1) );
13            free_value($1);
14            $1 = target;
15        }
16        $$ = $1;
17        yycode("%s += %s\n", target, rvalue( $3 ) );
18        free_value( $3 );
19    }
20    | e EQUAL e
21    {
22        if( *$1 != 'L' )
23            yyerror("Lvalue required\n");
24        else
25        {
26            if( *$3 != 'R' )           /* Convert $3 to an rvalue. */
27            {
28                yycode( "%s = %s\n", target = new_rvalue(), rvalue($3));
29                free_value($3);
30                $1 = target;
31            }
32            yycode("*%s = %s", *$1=='&' ? ($1)+1 : $1, $3); /* ★ */
33            $$ = $3;
34            free_value( $1 );
35        }
36    }
37    | ID
38    {
39        char buf[16];
40        sprintf( buf, "&%s", yytext );
41        /* ★ */
42        /* ★ */

```

Listing 6.66. continued...

```

41             $$ = strdup( buf );
42         }
43     ;
44 %% 
45 char *Lvalues[4] = { "L0", "L1", "L2", "L3" };
46 char *Lp       = Lvalues;
47
48 char *Rvalues[4] = { "R0", "R1", "R2", "R3" };
49 char *Rp       = Rvalues;
50
51 char *new_lvalue() { return *Lp++; } /* Allocate a new lvalue. */
52 char *new_rvalue() { return *Rp++; } /* Allocate a new rvalue. */
53
54 free_value( val ) /* Free an lvalue or an rvalue. */
55 char *val;
56 {
57     if (*val == 'L') --Lp = val; /* Physical lvalue */
58     else if( *val == 'R' ) --Rp = val;
59     else free( val ); /* Logical lvalue */
60 }
61
62 char *rvalue( value ) /* If the argument is an lvalue, return a string */
63 char *value; /* that can be used in an expression to reference */
64 { /* the lvalue's contents. */
65     static char buf[16];
66
67     if( *lvalue == 'R' ) /* It's already an rvalue. */
68         return value;
69
70     if( *lvalue == '&'amp; ) /* Logical lvalue */
71         return value + 1; /* *&_x cancels. Just use the name. */
72
73     sprintf( buf, "*%s", lvalue );
74     return buf;
75 }
```

The new-and-improved code-generation actions translate the input $x=y+z$ as follows:

```

R0 = _y ;
R0 += _z ;
_x = R0;
```

Summary: lvalues and rvalues.

To summarize all of the foregoing and put it into the context of a real compiler: two types of temporaries can be created as an expression is evaluated: an *lvalue* which holds the address of the object that, in turn, holds the result of the previous operation, and an *rvalue* which holds the result itself. The terms come from an expression like:

```
x = y;
```

The *rvalue*, *y* (on the right of the equal sign), is treated differently from the *lvalue*, *x* (on the left). The *y* must evaluate to its contents, but the *x* must evaluate to an address—to the place where the *y* is to be stored.

An expression that is said to “generate an *lvalue*”, can be seen as creating a temporary variable that holds an address of some sort, and the operator controls whether an *lvalue* or *rvalue* is generated. In the current C compiler, all identifier references generate *lvalues*—a temporary is created that holds the address of the referenced object. The

following operators also generate lvalues:

* [] . ->

Note that none of the foregoing operators do arithmetic—they all do a reference of some sort. Operators not in the foregoing list generate rvalues—a temporary is created that holds the result of the operation. In addition, the following operators require their operands to be lvalues:

++ -- = += -= etc.

Other operators can handle both lvalue and rvalue operands, but to use them, they must convert the lvalues to rvalues by removing one level of indirection.

A quick way to determine whether an expression forms an lvalue or an rvalue is to put the expression to the left of an equals sign and see if it's legal. Given the following declarations:

```
int x;
int array[10];
int *p = array;

struct tag
{
    int f1;
    int f2[5];
}
structure;
```

the following expressions generate lvalues:

x					
array	*	array		array[n]	
p	*	p	*	(p + n)	
structure.f1		structure.f2[n]	*	(structure.f2)	
				structure	(only if structure assignment is supported)

Note that both p and *p form lvalues, though the associated temporaries are of different types. The following are not lvalues because they can't appear to the left of an equal sign:

&x					
array					
structure.f2					
structure					(only if structure assignment is not supported)

6.8.3 Implementing Values, Higher-Level Temporary-Variable Support

This section describes the current compiler's implementation of lvalues and rvalues. The value structure, in Listing 6.67, is used for both of them, and all of the expression-processing productions pass around pointers to value structures as attributes. The various fields are used as follows:

- lvalue** True if the structure represents an lvalue, false if it's an rvalue.
- is_tmp** True if the structure represents a real temporary variable (an rvalue or a physical lvalue), false otherwise.
- offset** Position of the temporary variable in the stack frame. This number is the absolute value of the offset, in stack elements, from the base address of the temporary-variable region to the variable itself [i.e. the number that is returned from `tmp_alloc()`].

The value structure.

name	If the current value is an rvalue, this field is a string which, when used as an operand, evaluates to the value of the previously computed subexpression. If it's an lvalue, this field is a string that evaluates to the address of the object generated from the previous subexpression.
type etype	Points at the start and end of a type chain representing the value's type. If an rvalue, the name field references an object of the indicated type; if an lvalue, the name field evaluates to a pointer to an object of the indicated type. When an lvalue is created from a symbol (when an identifier is processed), a copy of the symbol's type chain is made and the copy is used in the value.
sym	When the lvalue is created from a symbol, this field points at that symbol. It's used only to print an occasional comment in the output, and is undefined if the lvalue was created by an operator like * or [].

Listing 6.67. *value.h*—Lvalues and Rvalues

```

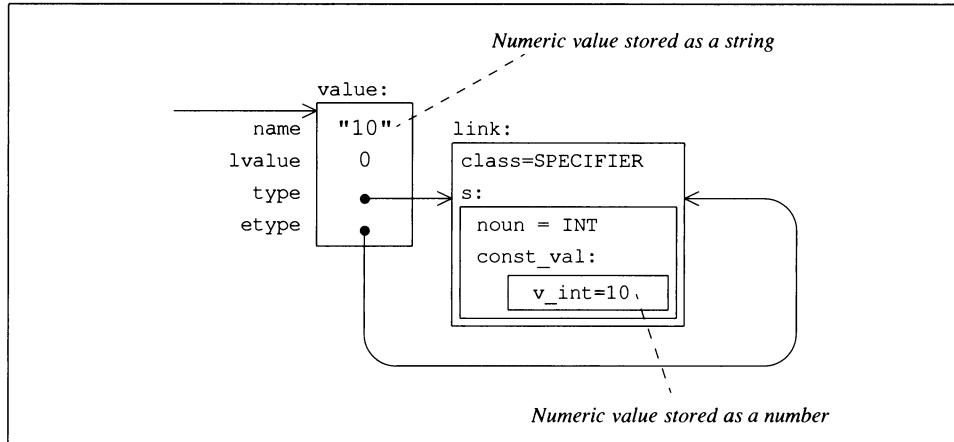
1  /* VALUE.H:      Various definitions for (l|r)values. "syntab.h" must be
2   *                  #included before #including this file.
3   */
4
5  #define VALNAME_MAX (NAME_MAX * 2)      /* Max. length of string in value.name */
6
7  typedef struct value
8  {
9    char     name[ VALNAME_MAX ]; /* Operand that accesses the value. */
10   link     *type;             /* Variable's type (start of chain). */
11   link     *etype;            /* Variable's type (end of chain). */
12   symbol   *sym;             /* Original symbol. */
13   unsigned lvalue   :1;       /* 1 = lvalue, 0 = rvalue. */
14   unsigned is_tmp  :1;       /* 1 if a temporary variable. */
15   unsigned offset  :14;      /* Absolute value of offset from base of */
16                           /* temporary region on stack to variable. */
17 } value;
18
19 #define LEFT  1           /* Second argument to shift_name() in value.c, */
20 #define RIGHT 0           /*          discussed below. */
21
22 #define CONST_STR(p) tconst_str((p)->type) /* Simplify tconst_str() calls */
23                                         /* with value arguments by */
24                                         /* extracting the type field. */

```

Constant value.

Note that values are used for constants as well as temporary variables. (There's a picture of one in Figure 6.19.) The attribute generated by *unop*→ICON is a pointer to a value structure for an rvalue. The type chain identifies the value as an **int** constant. [IS_CONSTANT(etype) in *syntab.h*, which examines the storage class of the right-most link in the chain, evaluates true.] Similarly, the numeric value of the constant is stored in the last link in the type chain (in the const_val union of the specifier substructure). The value structure's name field holds an ASCII string representing the number. The *unop*→NAME action that we'll look at momentarily creates an identical attribute if the identifier references an element of an enumerated type.

The advantage of this approach—using values for both constants and variables—is that all expression elements, regardless of what they are, synthesize value pointers as attributes.

Figure 6.19. A value that Represents an `int` Constant

Listing 6.68 contains various subroutines that manipulate value structures. The low-level allocation and deletion routines (`discard_value()` on line 46 of Listing 6.68 and `new_value()` on line 84 of the same listing) work just like the link-management routines discussed earlier. The values are allocated from `malloc()`, but discarded by placing them on the linked free list pointed to by `Value_free`, declared on line 14. Several structures are allocated at once—the number of structures is controlled by `VCHUNK`, defined on line 19. The type chain associated with the value is discarded at the same time as the value itself.

`new_value()`,
`discard_value()`

Listing 6.68. `value.c`—Low-Level value Maintenance

```

1  #include <stdio.h>
2  #include <tools/debug.h>
3  #include <tools/hash.h>
4  #include <tools/l.h>
5  #include <tools/c-code.h>
6  #include <tools/compiler.h>
7  #include "symtab.h"
8  #include "value.h"
9  #include "proto.h"
10 #include "label.h"
11
12 /*  VALUE.C:      Routines to manipulate lvalues and rvalues. */
13
14 PRIVATE value *Value_free = NULL;
15 #define VCHUNK 8           /* Allocate this many structure at a time. */
16
17 /*****
18
19 value  *new_value()
20 {
21     value *p;
22     int   i;
23
24     if( !Value_free )
25     {
26         if( !(Value_free = (value *) malloc( sizeof(value) * VCHUNK ) ) )
27         {
28             yyerror("INTERNAL, new_value: Out of memory\n");
29             exit( 1 );

```



Listing 6.68. continued...

```

30         }
31
32     for( p = Value_free, i = VCHUNK; --i > 0; ++p ) /* Executes VCHUNK-1 */
33         p->type = (link *) ( p+1 ); /* times. */
34
35     p->type = NULL ;
36 }
37
38     p = Value_free;
39     Value_free = (value *) Value_free->type;
40     memset( p, 0, sizeof(value) );
41     return p;
42 }
43
44 /*-----*/
45
46 void    discard_value( p )
47 value   *p;                                /* Discard a value structure */
48 {                                         /* and any associated links. */
49     if( p )
50     {
51         if( p->type )
52             discard_link_chain( p->type );
53
54         p->type = (link *)Value_free ;
55         Value_free = p;
56     }
57 }
58
59 /*-----*/
60
61 char    *shift_name( val, left )
62 value   *val;
63 int     left;
64 {
65     /* Shift the name one character to the left or right. The string might get
66      * truncated on a right shift. Returns pointer to shifted string.
67      * Shifts an & into the new slot on a right shift.
68      */
69
70     if( left )
71         memmove( val->name, val->name+1, sizeof(val->name) - 1 );
72     else
73     {
74         memmove( val->name+1, val->name, sizeof(val->name) - 1 );
75         val->name[ sizeof(val->name) - 1 ] = '\0' ;
76         val->name[0] = '&';
77     }
78
79     return val->name;
80 }
81
82 /*-----*/
83
84 char    *rvalue( val )
85 value   *val;
86 {
87     /* If val is an rvalue, do nothing and return val->name. If it's an lvalue,
88      * convert it to an rvalue and return the string necessary to access the
89      * rvalue. If it's an rvalue constant, return a string representing the
90      */

```

Listing 6.68. continued...

```

90     * number (and also put that string into val->name).
91     */
92
93     static char buf[ NAME_MAX + 2 ] = "*" ;
94
95     if( !val->lvalue )                                /* It's an rvalue already. */
96     {
97         if( IS_CONSTANT(val->etype) )
98             strcpy( val->name, CONST_STR(val) );
99         return val->name;
100    }
101   else
102   {
103       val->lvalue = 0;
104       if( * val->name == '&'amp; )                      /* *& cancels */
105       {
106           shift_name(val, LEFT);
107           return val->name;
108       }
109   else
110   {
111       strncpy( &buf[1], val->name, NAME_MAX+1 );
112       return buf;
113   }
114 }
115 }
116
117 char *rvalue_name( val )
118 value *val;
119 {
120     /* Returns a copy of the rvalue name without actually converting val to an
121      * rvalue. The copy remains valid only until the next rvalue_name call.
122      */
123
124     static char buf[ (NAME_MAX * 2) + 1 ];           /* sizeof(val->name) + 1 */
125
126     if( !val->lvalue )
127     {
128         if( IS_CONSTANT(val->etype) ) strcpy( buf, CONST_STR(val) );
129         else                           strcpy( buf, val->name );
130     }
131   else
132   {
133       if( * val->name == '&'amp; ) strcpy ( buf,        val->name + 1 );
134       else                           sprintf( buf, "*%s", val->name );
135   }
136   return buf;
137 }
```

The `rvalue()` subroutine on line 84 of Listing 6.68 is used heavily in the expression-processing code. It is used to fetch the actual value associated with either an `rvalue` or `lvalue`. It is passed a `value` structure. If this value is an `rvalue` that represents a numeric constant, a string representing the constant's numeric value is returned. If it's a normal `rvalue`, the `name` field is returned.

If the argument is a pointer to an `lvalue`, the `value` is converted to an `rvalue` and a string that can be used to access the `rvalue` is returned. The conversion to `rvalue` involves modifying two fields. First, the `lvalue` field is set false. Then, if this is a

`rvalue()`
`rvalue()` and `rvalues`.

`rvalue()` and logical
`lvalues`.

`rvalue()` and physical lvalues.

logical lvalue, the name field is shifted one character to the left and the resulting string is returned.

If the incoming value represents a physical lvalue, a string containing the name preceded by a * is returned, but the name field isn't modified. This behavior is actually an anomaly because a physical rvalue is not created. That is, to be consistent with the logical lvalue case, the `rvalue()` subroutine should emit code to copy the object pointed to by the lvalue into a physical temporary variable. It should then modify the value structure to represent this new temporary rather than the original one. I haven't done this here because I wanted the code-generation actions to be able to control the target. For example, the following is taken from the code-generation action for a return statement. \$2 is a pointer to a value for the object being returned. If the input takes the form `return x`, then \$2 points at an lvalue. If the input takes the form `return x+1`, then \$2 points at an rvalue:

```
gen( " =", IS_INT($2->type) ? "rF.w.low" : "rF.l", rvalue($2) );
```

The target (the physical rvalue) has to be one of two registers, not a temporary variable. Other situations arise where the target is a temporary, so I simplified the interface to `rvalue()` by requiring the caller to generate the physical rvalue if necessary. The string that's returned from `rvalue()` can always be used as the source operand in the assignment. The `rvalue_name()` subroutine on line 117 returns the same thing as `rvalue()`, but it doesn't modify the value structure at all, it just returns a string that you can use as a source operand.

`rvalue_name()`

Temporary-variable
`value_tmpl_create()`.

`get_prefix()`

The value-maintenance functions continue in Listing 6.69 with a second layer of higher-level functions. Temporary-variable values are also handled here. Two routines, starting on line 138, are provided for temporary-variable creation. `tmp_create()` is the lower-level routine. It is passed a type and creates an rvalue for that type using the low-level allocation routine discussed earlier [`tmp_alloc()`] to get space for the variable on the stack. A copy of the input type string is usually put into the value. An `int` rvalue is created from scratch if the type argument is `NULL`, however. If the second argument is true, a link for a pointer declarator is added to the left of the value's type chain. The value's name field is created by `tmp_create()`. It initialized to a string which is output as the operand of an instruction, and which accesses the temporary variable. The name is generated on line 175, and it looks something like this: `WP(T(0))`. The type component (`WP`) changes with the actual type of the temporary. It's created by `get_prefix()` on line 181. The `0` is the offset from the base of the temporary-variable region to the current variable. We saw the `T()` macro earlier. It translates the offset to a frame-pointer reference.

Listing 6.69. `value.c`—High-Level value Maintenance

```

138  value *tmp_create( type, add_pointer )
139  link  *type;           /* Type of temporary, NULL to create an int.          */
140  int   add_pointer;    /* If true, add pointer declarator to front of type */
141  {                   /* before creating the temporary.                      */
142
143      /* Create a temporary variable and return a pointer to it. It's an rvalue
144      * by default.
145      */
146
147      value *val;
148      link  *lp;
149

```



Listing 6.69. continued...

```

150     val           = new_value();
151     val->is_tmp = 1;
152
153     if( type )                                /* Copy existing type. */
154         val->type   = clone_type( type, &lp );
155     else                                         /* Make an integer.      */
156     {
157         lp           = new_link();
158         lp->class   = SPECIFIER;
159         lp->NOUN    = INT;
160         val->type   = lp;
161     }
162
163     val->etype   = lp;
164     lp->SCLASS   = AUTO;                      /* It's an auto now, regardless      */
165                                         /* of the previous storage class. */
166     if( add_pointer )
167     {
168         lp           = new_link();
169         lp->DCL_TYPE = POINTER;
170         lp->next     = val->type;
171         val->type   = lp;
172     }
173
174     val->offset = tmp_alloc( get_size( val->type ) );
175     sprintf(val->name, "%s( T(%d) )", get_prefix(val->type), (val->offset + 1));
176     return( val );
177 }
178
179 /*-----*/
180
181 char    *get_prefix( type )
182 link    *type;
183 {
184     /* Return the first character of the LP(), BP(), WP(), etc., directive
185      * that accesses a variable of the given type. Note that an array or
186      * structure type is assumed to be a pointer to the first cell at run time.
187      */
188
189     int   c;
190
191     if( type )
192     {
193         if( type->class == DECLARATOR )
194         {
195             switch( type->DCL_TYPE )
196             {
197                 case ARRAY:   return( get_prefix( type->next ) );
198                 case FUNCTION: return PTR_PREFIX;
199                 case POINTER:
200                     c = *get_prefix( type->next );
201
202                     if       ( c == *BYTE_PREFIX ) return BYTEPTR_PREFIX;
203                     else if( c == *WORD_PREFIX ) return WORDPTR_PREFIX;
204                     else if( c == *LWORD_PREFIX ) return LWORDPTR_PREFIX;
205
206                     break;
207             }
208         }
209     }

```



Listing 6.69. continued...

```

209         else
210     {
211         switch( type->NOUN )
212         {
213             case INT:      return (type->LONG) ? LWORD_PREFIX : WORD_PREFIX;
214             case CHAR:
215             case STRUCTURE: return BYTEPTR_PREFIX ;
216         }
217     }
218 }
219 yyerror("INTERNAL, get_prefix: Can't process type %s.\n", type_str(type) );
220 exit( 1 );
221 }
222
223 /*-----*/
224
225 value *tmp_gen( tmp_type, src )
226 link   *tmp_type;          /* type of temporary taken from here           */
227 value   *src;            /* source variable that is copied into temporary */
228 {
229     /* Create a temporary variable of the indicated type and then generate
230      * code to copy src into it. Return a pointer to a "value" for the temporary
231      * variable. Truncation is done silently; you may want to add a lint-style
232      * warning message in this situation, however. Src is converted to an
233      * rvalue if necessary, and is released after being copied. If tmp_type
234      * is NULL, an int is created.
235 */
236
237 value   *val;                  /* temporary variable                         */
238 char    *reg;
239
240 if ( !the_same_type( tmp_type, src->type, 1 ) )
241 {
242     /* convert_type() copies src to a register and does any necessary type
243      * conversion. It returns a string that can be used to access the
244      * register. Once the src has been copied, it can be released, and
245      * a new temporary (this time of the new type) is created and
246      * initialized from the register.
247 */
248
249     reg = convert_type( tmp_type, src );
250     release_value( src );
251     val = tmp_create( IS_CHAR(tmp_type) ? NULL : tmp_type, 0 );
252     gen( "=", val->name, reg );
253 }
254 else
255 {
256     val = tmp_create( tmp_type, 0 );
257     gen( "=", val->name, rvalue(src) );
258     release_value( src );
259 }
260 return val;
261 }
262
263 /*-----*/
264
265 char    *convert_type( targ_type, src )
266 link   *targ_type;          /* type of target object                      */
267 value   *src;              /* source to be converted                     */
268 {

```

Listing 6.69. continued...

```

269     int      dsiz;                      /* src size, dst size          */
270     static char reg[16];                /* place to assemble register name */
271
272     /* This routine should only be called if the target type (in targ_type)
273      * and the source type (in src->type) disagree. It generates code to copy
274      * the source into a register and do any type conversion. It returns a
275      * string that references the register. This string should be used
276      * immediately as the source of an assignment instruction, like this:
277      *
278      *     gen( "=", dst_name, convert_type( dst_type, src ); )
279      */
280
281     sprintf( reg, "r0.%s", get_suffix(src->type) );           /* copy into */
282     gen( "=", reg, rvalue(src) );                                /* register. */
283
284     if( (dsiz = get_size(targ_type)) > get_size(src->type) )
285     {
286         if( src->etype->UNSIGNED )                               /* zero fill */
287         {
288             if( dsiz == 2 ) gen( "=", "r0.b.bl", "0" );
289             if( dsiz == 4 ) gen( "=", "r0.w.high", "0" );
290         }
291     else
292         {
293             if( dsiz == 2 ) gen( "ext_low", "r0" );
294             if( dsiz == 4 ) gen( "ext_word", "r0" );
295         }
296     }
297
298     sprintf( reg, "r0.%s", get_suffix(targ_type) );
299     return reg;
300 }
301
302 /*-----*/
303
304 PUBLIC int      get_size( type )
305 link   *type;
306 {
307     /* Return the number of bytes required to hold the thing referenced by
308      * get_prefix().
309      */
310
311     if( type )
312     {
313         if( type->class == DECLARATOR )
314             return (type->DCL_TYPE == ARRAY) ? get_size(type->next) : PSIZE ;
315         else
316         {
317             switch( type->NOUN )
318             {
319                 case INT:      return (type->LONG) ? LSIZE : ISIZE;
320                 case CHAR:
321                 case STRUCTURE: return CSIZE;
322             }
323         }
324     }
325     yyerror("INTERNAL, get_size: Can't size type: %s.\n", type_str(type));
326     exit(1);
327 }
328

```



Listing 6.69. continued...

```

329  /*-----*/
330
331 char      *get_suffix( type )
332 link      *type;
333 {
334     /* Returns the string for an operand that gets a number of the indicated
335      * type out of a register. (It returns the xx in : r0.xx). "pp" is returned
336      * for pointers, arrays, and structures--get_suffix() is used only for
337      * temporary-variable manipulation, not declarations. If an array or
338      * structure declarator is a component of a temporary-variable's type chain,
339      * then that declarator actually represents a pointer at run time. The
340      * returned string is, at most, five characters long.
341
342
343     if( type )
344     {
345         if( type->class == DECLARATOR )
346             return "pp";
347         else
348             switch( type->NOUN )
349             {
350                 case INT:       return (type->LONG) ? "l" : "w.low";
351                 case CHAR:
352                 case STRUCTURE: return "b.b0";
353             }
354
355     yyerror("INTERNAL, get_suffix: Can't process type %s.\n", type_str(type) );
356     exit( 1 );
357 }
358
359 /*-----*/
360
361 void      release_value( val )
362 value    *val;           /* Discard a value, first freeing any space */
363 {                   /* used for an associated temporary variable. */
364     if( val )
365     {
366         if( val->is_tmp )
367             tmp_free( val->offset );
368         discard_value( val );
369     }
370 }
```

Create and initialize temporary: `tmp_gen()`.

`tmp_gen()` on line 225 of Listing 6.69 both creates the temporary and emits the code necessary to initialize it. (`tmp_create()` doesn't initialize the temporary.) The subroutine is passed three arguments. The first is a pointer to a type chain representing the type of the temporary. An `int` is created if this argument is `NULL`. The second argument is a pointer to a `value` representing an object to be copied into the temporary. Code is generated to do the copying, and the source `value` is released after the code to do the copy is generated. If the source is an lvalue, it is converted to an rvalue first.

Type conversions.

`convert_type()`.

If the source variable and temporary variable's types don't match, code is generated to do any necessary type conversions. Type conversions are done in `convert_type()` on line 265 of Listing 6.69 by copying the original variable into a register; if the new variable is larger than the source, code to do sign extension or zero fill is emitted. The subroutine returns a string that holds the name of the target register. For example, the input code:

```
int i;
long l;
foo() { i=i+1; }
```

generates the following output:

```
r0.w.low = W(&_i);
ext_word(r0);
L( T(1) ) = r0.l;

L( T(1) ) += L(&_l);

r0.l      = L( T(1) );
W(&_i)    = r0.w.low;
```

The first two lines, emitted by `convert_type()`, create a **long** temporary variable and initialize it from an **int**. The code copies *i* into a word register and then sign extends the register to form an lword. The string "r0.l" is passed back to `tmp_gen()`, which uses it to emit the assignment on the third line. The last two lines demonstrate how a truncation is done. `convert_type` emits code to copy `T(1)` into a register. It then passes the string "r0.w.low" back up to the calling routine, which uses it to generate the assignment.

The test on line 284 of Listing 6.69 checks to see if sign extension is necessary (by comparing the number of bytes used for both variables—`get_size()` starts on line 304 of Listing 6.69). Code to do the type conversion is done on lines 286 to 295 of Listing 6.69. `get_suffix`, used on line 298 to access the register, starts on 331 of listing 6.69.

The `release_value()` subroutine on line 361 of Listing 6.69 is a somewhat higher-level version of `discard_value()`. It's used for temporary variables and both discards the value structure and frees any stack space used for the temporary. Note that the original source value is discarded on line 250 of Listing 6.69 as soon as the variable is copied into a register, so the memory used by the original variable can be recycled for the new temporary.

`release_value()`.

6.8.4 Unary Operators

Armed with the foregoing, you can proceed to tackle expression processing, starting with the unary operators and working our way up the grammar to complete expressions. Because you're starting at the bottom, it is helpful to look at the overall structure of the expression productions in Appendix C before continuing. Start with the *expr* nonterminal and go to the end of the grammar.

The unary operators are handled by the right-hand sides of the *unary* nonterminal, which you get to like this:

<i>compound_stmt</i>	→	LC local_defs stmt_list RC
<i>stmt_list</i>	→	<i>stmt</i>
<i>stmt</i>	→	<i>expr SEMI</i>
<i>expr</i>	→	<i>binary</i> <i>unary</i>

(I've both simplified and left out a few intermediate productions for clarity. See Appendix C for the actual productions.) The simplest of the *unary* productions are in Listing 6.70. The top production just handles parenthesized subexpressions, and the second right-hand side recognizes, but ignores, floating-point constants.

The right-hand side on line 597 handles integer constants, creating a value structure like the one examined earlier in Figure 6.19. `make_icon()`, in Listing 6.71, does the actual work, creating a value for the integer constant, which is then passed back up as an attribute. This subroutine is used in several other places so it does more than is

Integer constants,
`make_icon()`.

Listing 6.70. *c.y*— Unary Operators (Part 1)—Constants and Identifiers

```

594 unary
595     : LP expr RP      { $$ = $2; }
596     | FCON           { yyerror("Floating-point not supported\n"); }
597     | ICON            { $$ = make_icon ( yytext, 0 ); }
598     | NAME            { $$ = do_name   ( yytext, $1 ); }

```

required by the current production. In particular, the incoming parameter can represent the number either as an integer or as a string. Numeric input is used here by setting *yytext* to NULL.

Identifiers,
unary→**NAME**.

do_name().

Identifiers are handled by on line 598 of Listing 6.70, above. The attribute attached to the **NAME** at \$1 is a pointer to the symbol table entry for the identifier or NULL if the identifier isn't in the table. The actual work is done by *do_name()*, in Listing 6.72, below.

Listing 6.71. *value.c*— Make Integer and Integer-Constant Rvalues

```

371 value *make_icon( yytext, numeric_val )
372 char *yytext;
373 int numeric_val;
374 {
375     /* Make an integer-constant rvalue. If yytext is NULL, then numeric_val
376     * holds the numeric value, otherwise the second argument is not used and
377     * the string in yytext represents the constant.
378     */
379
380     value *vp;
381     link *lp;
382     char *p;
383
384     vp      = make_int();
385     lp      = vp->type;
386     lp->SCLASS = CONSTANT;
387
388     if( !yytext )
389         lp->V_INT = numeric_val;
390
391     else if( *yytext == '\'' )
392     {
393         ++yytext;                  /* Skip the quote. */
394         lp->V_INT = esc( &yytext );
395     }
396     else                         /* Initialize the const_val field      */
397     {                           /* based on the input type. stoul()    */
398                               /* converts a string to unsigned long.*/
399         for( p = yytext; *p ; ++p )
400         {
401             if      ( *p=='u' || *p=='U' ) { lp->UNSIGNED = 1; }
402             else if( *p=='l' || *p=='L' ) { lp->LONG      = 1; }
403         }
404         if( lp->LONG )
405             lp->V ULONG = stoul( &yytext );
406         else
407             lp->V_UINT  = (unsigned int) stoul( &yytext );
408     }
409     return vp;
410 }

```



Listing 6.71. continued...

```

411
412 /*-----*/
413
414 value *make_int()
415 {
416     link *lp;
417     value *vp;
418
419     lp      = new_link();
420     lp->class = SPECIFIER;
421     lp->NOUN  = INT;
422     vp      = new_value();           /* It's an rvalue by default. */
423     vp->type = vp->etype = lp;
424     return vp;
425 }
```

Listing 6.72. op.c—Identifier Processing

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/debug.h>
4 #include <tools/hash.h>
5 #include <tools/l.h>
6 #include <tools/compiler.h>
7 #include <tools/c-code.h>
8 #include "symtab.h"
9 #include "value.h"
10 #include "proto.h"
11 #include "label.h"
12
13 /* OP.C      This file contains support subroutines for the arithmetic */
14 /*          operations in c.y. */
15
16 symbol *Undecl = NULL; /* When an undeclared symbol is used in an expression,
17                         * it is added to the symbol table to suppress subse-
18                         * quent error messages. This is a pointer to the head
19                         * of a linked list of such undeclared symbols. It is
20                         * purged by purge_undecl() at the end of compound
21                         * statements, at which time error messages are also
22                         * generated.
23 */
24
25 /*-----*/
26
27 value *do_name( yytext, sym )
28 char *yytext;             /* Lexeme
29 symbol *sym;              /* Symbol-table entry for id, NULL if none. */
30 {
31     link *chain_end, *lp ;
32     value *synth;
33     char buf[ VALNAME_MAX ];
34
35 /* This routine usually returns a logical lvalue for the referenced symbol.
36 * The symbol's type chain is copied into the value and value->name is a
37 * string, which when used as an operand, evaluates to the address of the
38 * object. Exceptions are aggregate types (arrays and structures), which
39 * generate pointer temporary variables, initialize the temporaries to point
```

Listing 6.72. continued...

```

40      * at the first element of the aggregate, and return an rvalue that
41      * references that pointer. The type chain is still just copied from the
42      * source symbol, so a structure or pointer must be interpreted as a pointer
43      * to the first element/field by the other code-generation subroutines.
44      * It's also important that the indirection processing (* [] . ->) set up
45      * the same sort of object when the referenced object is an aggregate.
46      *
47      * Note that !sym must be checked twice, below. The problem is the same one
48      * we had earlier with typedefs. A statement like foo(){int x;x=1;} fails
49      * because the second x is read when the semicolon is shifted--before
50      * putting x into the symbol table. You can't use the trick used earlier
51      * because def_list is used all over the place, so the symbol table entries
52      * can't be made until local_defs->def_list is processed. The simplest
53      * solution is just to call findsym() if NULL was returned from the scanner.
54      *
55      * The second if(!sym) is needed when the symbol really isn't there.
56      */
57
58  if( !sym )
59      sym = (symbol *) findsym( Symbol_tab, yytext );
60
61  if( !sym )
62      sym = make_implicit_declaration( yytext, &Undecl );
63
64  if( IS_CONSTANT(sym->etype) )           /* it's an enum member */
65  {
66      if( IS_INT(sym->type) )
67          synth = make_icon( NULL, sym->type->V_INT );
68      else
69      {
70          yyerror("Unexpected noninteger constant\n");
71          synth = make_icon( NULL, 0 );
72      }
73  }
74  else
75  {
76      gen_comment("%s", sym->name);    /* Next instruction will have symbol */
77      /* name as a comment. */           */
78
79  if( !(lp = clone_type( sym->type, &chain_end)) )
80  {
81      yyerror("INTERNAL do_name: Bad type chain\n" );
82      synth = make_icon( NULL, 0 );
83  }
84  else if( IS_AGGREGATE(sym->type) )
85  {
86      /* Manufacture pointer to first element */
87
88      sprintf(buf, "&%s(%s)",
89                  IS_ARRAY(sym->type) ? get_prefix(lp) : BYTE_PREFIX,
90                  sym->rname );
91
92      synth = tmp_create(sym->type, 0);
93      gen( "=" , synth->name, buf );
94  }
95  else
96  {
97      synth      = new_value();
98      synth->lvalue = 1    ;
99      synth->type   = lp   ;

```



Listing 6.72. continued...

```

100         synth->etype = chain_end;
101         synth->sym = sym ;
102
103     if( sym->implicit || IS_FUNCT(lp) )
104         strcpy( synth->name, sym->rname );
105     else
106         sprintf(synth->name,
107             (chain_end->SCLASS == FIXED) ? "&%s(&%s)" : "&%s(%s)",
108             get_prefix(lp), sym->rname);
109     }
110 }
111
112 return synth;
113 }
114
115 /*-----*/
116
117 symbol *make_implicit_declaration( name, undeclp )
118 char *name;
119 symbol **undeclp;
120 {
121     /* Create a symbol for the name, put it into the symbol table, and add it
122      * to the linked list pointed to by *undeclp. The symbol is an int. The
123      * level field is used for the line number.
124 */
125
126     symbol *sym;
127     link *lp;
128     extern int yylineno;      /* created by LeX */
129     extern char *yytext;
130
131     lp = new_link();
132     lp->class = SPECIFIER;
133     lp->NOUN = INT;
134     sym = new_symbol( name, 0 );
135     sym->implicit = 1;
136     sym->type = sym->etype = lp;
137     sym->level = yylineno;    /* Use the line number for the declaration */
138                             /* level so that you can get at it later */
139                             /* if an error message is printed. */
140
141     sprintf( sym->rname, "_%1.*s", sizeof(sym->rname)-2, yytext );
142     addsym ( Symbol_tab, sym );
143
144     sym->next = *undeclp;           /* Link into undeclared list. */
145     *undeclp = sym;
146
147     return sym;
148 }
149
150 /*-----*/
151
152 PUBLIC void purge_undecl()
153 {
154     /* Go through the undeclared list. If something is a function, leave it in
155      * the symbol table as an implicit declaration, otherwise print an error
156      * message and remove it from the table. This routine is called at the
157      * end of every subroutine.
158 */
159

```



Listing 6.72. continued...

```

160     symbol *sym, *cur ;
161
162     for( sym = Undecl; sym; )
163     {
164         cur      = sym;           /* remove current symbol from list */
165         sym      = sym->next;
166         cur->next = NULL;
167
168         if( cur->implicit )
169         {
170             yyerror("%s (used on line %d) undeclared\n", cur->name, cur->level);
171             delsym( Symbol_tab, cur );
172             discard_symbol( cur );
173         }
174     }
175     Undecl = NULL;
176 }
```

The incoming `sym` argument is `NULL` when the scanner can't find the symbol in the table. The table lookup on line 58 of Listing 6.72 takes care of the time-delay problem discussed earlier when types were discussed. Code like this:

```

int x;
x = 1;
```

won't work because the symbol-table entry for `x` isn't created until after the second `x` is read. Remember, the action is performed after the **SEMI** is shifted, and the lookahead character is read as part of the shift operation.

Undeclared identifiers.

Undeclared identifiers pose a particular problem. The difficulty is function calls: an implicit declaration for the function must be created the first time the function is used. Undeclared variables are hard errors, though. Unfortunately, there's no way for the **unop→NAME** production to know how the name is being used. It doesn't know whether the name is part of a function call or not.

make_implicit_declaration()

If the symbol really isn't in the table, `do_name()` creates an implicit declaration of type `int` for the undeclared identifier (on line 62 of Listing 6.72). The `make_implicit_declaration()` subroutine is found on lines 117 to 148 of Listing 6.72. The new symbol is put into the symbol table, and the cross links form a linked list of undeclared symbols. The head-of-list pointer is `Undecl`, declared on line 16. The implicit symbol is marked as such by setting the `implicit` bit true on line 137 of Listing 6.72. The implicit symbol is modified when a function call is processed. A "function" declarator is added to the front of the type chain, and the `implicit` bit is turned off. I'll discuss this process in detail in a moment.

purge_undecl()

After the entire subroutine is processed and the subroutine suffix is generated, `purge_undecl()` (on line 152 of Listing 6.72) is called to find and delete the undeclared symbols. This subroutine traverses the list of implicitly-generated symbols and prints "undeclared symbol" error messages for any of them that aren't functions. It also removes the undeclared variables from the symbol table and frees the memory.

Symbolic constants.

One advantage of this approach is that the "undeclared symbol" error message is printed only once, regardless of the number of times that the symbol is used. The disadvantage is that the error message isn't printed until the entire subroutine has been processed.

Returning to the normal situation, by the time you get to line 64 of Listing 6.72, `sym` points at a symbol for the current identifier. If that symbol represents a constant, then it

was created by a previous enumerated-type declaration. The action in this case (on lines 66 to 72) creates an rvalue for the constant, just as if the number, rather than the symbolic name, had been found.

If the current symbol is an identifier, a value is created on lines 74 to 110 of Listing 6.72. The type chain is cloned. Then, if the symbol represents an aggregate type such as an array, a temporary variable is generated and initialized to point at the first element of the aggregate. Note that this temporary is an rvalue, not an lvalue. (An array name without a star or brackets is illegal to the left of an equal sign.) It's also important to note that arrays and pointers can be treated identically by subsequent actions because a physical pointer to the first array element is generated—pointers and arrays are physically the same thing at run time. A **POINTER** declarator and an **ARRAY** declarator are treated identically by all of the code that follows. I'll explore this issue further, along with a few examples, in a moment.

If the current symbol is not an aggregate, a logical lvalue is created in the **else** clause on lines 95 to 109 of Listing 6.72. The value's name field is initialized to a string that, when used as an operand in a C-code instruction, evaluates either to the address of the current object or to the address of the first element if it's an aggregate.

If the symbol is a function or if it was created with an implicit declaration, the name is used as is. (We're on lines 103 of Listing 6.72.) I'm assuming that an implicit symbol will eventually reference a function. The compiler generates bad code if the symbol doesn't, but so what? It's a hard error to use an undeclared variable.

All other symbols are handled by the `sprintf()` call on line 106 of Listing 6.72. The strategy is to represent all possible symbols in the same way, so that subsequent actions won't have to worry about how the name is specified. In other words, a subsequent action shouldn't have to know whether the referenced variable is on the stack, at a fixed memory address, or whatever. It should be able to use the name without thinking.

You can get this consistency with one of the stranger-looking, C-code storage classes. The one case that gives us no flexibility is a frame-pointer-relative variable, which must be accessed like this: `WP(fp+8)`. This string starts with a type indicator (`WP`), the variable reference itself is an address (`fp+8`), and an indirect addressing mode must be used to access the variable. To be consistent, variables at fixed addresses such as global variables are represented in a similar way—like this: `WP(&_p)`. The `fp+8` and `_p` are just taken from the symbol's name field. The type (`WP`, here) is figured by `get_prefix()` by looking at the variable's actual type. Since we are forming lvalues, the names must evaluate to addresses, so we need one more ampersand in both cases. `&WP(fp+8)` and `&WP(&_p)` are actually put into the name array.

The next *unop* right-hand side, which handles string constants, is in Listing 6.73. The `string_const` productions in Listing 6.74 collect adjacent string constants and concatenate them into the `Str_buf` array, declared in the `occs` declarations section at the top of listing 6.74. The attribute attached to the `string_const` is a pointer to the assembled string, which is then turned into an rvalue by the `make_scon()` call on line 602 of Listing 6.73. The subroutine itself is in Listing 6.75. String constants are of type “pointer to `char`”. The string itself is not stored internally by the compiler beyond the time necessary to parse it. When a string constant is recognized, a declaration like the following one is output to the `data` segment:

```
private char S156[] = "contents of string";
```

The numeric component of the label (156) is remembered in `v_int` field of the last link in the type chain, and the entire label is also loaded into the value's name array.

The next operator, `sizeof`, is handled by the productions in Listing 6.76. Three cases have to be handled separately:

Creating identifier values.

Aggregate types represented by physical pointers at run time.

`value.name`

All symbols represented identically.

Using `&WP(&_p)` for lvalues

String constants.

The `sizeof` operator.

Listing 6.73. c.y— Unary Operators (Part 2)—String Constants

```

599 /* unop: */
600     | string_const %prec COMMA
601     {
602         $$ = make_scon( $1 );
603         yydata( "private\tchar\t%s[]=\"%s\";\n", $$->name,$1 );
604     }

```

Listing 6.74. c.y— Array in which String Constants are Assembled

```

139  %{
140  #define STR_MAX 512           /* Maximum size of a string constant. */
141  char   Str_buf[STR_MAX];    /* Place to assemble string constants. */
142  %}

605 string_const
606     : STRING
607     {
608         $$      = Str_buf;
609         *Str_buf = '\0';
610
611         yytext[ strlen(yytext) - 1 ] = '\0'; /* Remove trailing " */
612
613         if( concat(STR_MAX, Str_buf, Str_buf, yytext+1, NULL) < 0 )
614             yyerror("String truncated at %d characters\n", STR_MAX );
615     }
616     | string_const STRING
617     {
618         yytext[ strlen(yytext) - 1 ] = '\0'; /* Remove trailing " */
619
620         if( concat(STR_MAX, Str_buf, Str_buf, yytext+1, NULL) < 0 )
621             yyerror("String truncated at %d characters\n", STR_MAX );
622     }
623 ;

```

sizeof("string") evaluates to the number of characters in the string.
sizeof(expression) evaluates to the size of the temporary variable that holds the result of the evaluation.

```

int x;
long y;
sizeof( x ); /* = 2 (size of a word). */
sizeof( y ); /* = 4 (size of an lword). */
sizeof( x+y ); /* = 4 (size of an lword). */

```

sizeof(type) evaluates to the number of bytes required by a variable of the indicated type.

The returned attribute for all three cases is an integer-constant rvalue, the *const_val* field of which holds the required size. The attribute attached to the *expr* (\$3 in the second right-hand side in Listing 6.76) is a value for the temporary that holds the result of the run-time expression evaluation. Its *type* field determines the size. Similarly, the attribute attached to the *abstract_decl* (\$3 in the third right-hand side in Listing 6.76) is a nameless symbol that holds a type chain for the indicated type.

Listing 6.75. *value.c*— Make a String-Constant Rvalue

```

426 value *make_scon( str )
427 char *str;
428 {
429     link *p;
430     value *synth;
431     static unsigned label = 0;
432
433     synth = new_value();
434     synth->type = new_link();
435
436     p = synth->type;
437     p->DCL_TYPE = POINTER;
438     p->next = new_link();
439
440     p = p->next;
441     p->class = SPECIFIER;
442     p->NOUN = CHAR;
443     p->SCLASS = CONSTANT;
444     p->V_INT = ++label;
445
446     synth->etype = p;
447     sprintf( synth->name, "%s%d", L_STRING, label );
448     return synth;
449 }
```

Listing 6.76. *c.y*— Unary Operators (Part 3)—**sizeof**

```

624 /* unop: */
625     | SIZEOF LP string_const RP           %prec SIZEOF
626     { $$ = make_icon(NULL,strlen($3) + 1); }
627
628     | SIZEOF LP expr RP                %prec SIZEOF
629     { $$ = make_icon(NULL,get_sizeof($3->type));
630         release_value( $3 );
631     }
632     | SIZEOF LP abstract_decl RP      %prec SIZEOF
633     {
634         $$ = make_icon( NULL,
635                           get_sizeof($3->type));
636         discard_symbol( $3 );
637     }
```

So far, none of the operators have generated code. With Listing 6.77, we move on to an operator that actually does something. The right-hand side in Listing 6.77 processes the cast operator. The test on line 641 checks to see if the cast is legal. If it is, and the source and destination types are both pointers, it just replaces the type chain in the operand with the one from the cast operator on lines 648 to 651. Note that this behavior, though common in C compilers, can cause problems because the contents of the pointer are not changed, only the type. A run-time alignment error occurs the first time you try to use the pointer if the alignment isn't correct.

The third possibility, on line 655 of Listing 6.77, takes care of all other cases. The `tmp_gen()` call creates a temporary of the required type, copies the contents of the operand into the temporary (making any necessary type conversions and translating

Type-conversion operators.

Casts.

Listing 6.77. c.y— Unary Operators (Part 4)—Casts and Address-of

```

638 /* unop: */
639     | LP abstract_decl RP unary
640     {
641         if( IS_AGGREGATE($2->type) )
642         {
643             yyerror( "Illegal cast to aggregate type\n" );
644             $$ = $4;
645         }
646         else if( IS_POINTER($2->type) && IS_POINTER($4->type) )
647         {
648             discard_link_chain( $4->type );
649             $4->type = $2->type;
650             $4->etype = $2->type;
651             $$ = $4;
652         }
653         else
654         {
655             $$ = tmp_gen( $2->type, $4 );
656         }
657     }

```

lvalues to rvalues as necessary), and finally, it deletes the value associated with the operand.

Unary minus, NOT, one's complement.

The next issue is the unary arithmetic operators, handled by the right-hand sides in Listing 6.78. The first one takes care of unary minus, the second handles the logical NOT and one's-complement operators (`!`, `~`). The attribute attached to the UNOP is the lexeme (passed in an `int` because it's only one character). Again, the work is done by a subroutine: `do_unop()`, at the top of Listing 6.79.

The `~` and unary `-` operators.

The arithmetic operators are easy. The action (on lines 184 to 188 of Listing 6.79) just modifies the operand, passing the operator through to `gen()`. If the name field of the operand holds the string "`W(fp+8)`", then the following code is generated for unary minus and one's complement:

```

W(T(0)) =- W(fp+8); /* -name */
W(T(0)) =~ W(fp+8); /* ~name */

```

In the first case, the `-` is supplied in the production, in the second, the attribute is used.

The `!` operator.

The logical NOT operator is a little harder to do because it must evaluate either to 1 or 0, depending on the truth or falsity of the operand. Assuming that `x` is at `W(fp+6)`, the expression `!x` generates the following code:

```

EQ( W(fp+6), 0 )                                /* if the operand if false */
    goto T0;
F0:   W( T(0) ) = 0;                               /* temp = 0 */
    goto E0;
T0:   W( T(0) ) = 1;                               /* temp = 1 */
    E0:

```

`tf_label()`

The synthesized attribute is the temporary variable that holds the 1 or 0. All but the first two instructions are created by `gen_false_true()`, on line 206 of Listing 6.79. The T label takes care of the true situation, the F label with false, and the E (for *end*) lets the false case jump over the true one after making the assignment. The numeric component of all three labels is the same, and is generated by the `tf_label()` call on line 192 of Listing 6.79. (`tf_label()` is declared just a few lines down, on line 199.) This label is then passed to `gen_false_true` as an argument.

All but the first two lines of the output are emitted by `gen_false_true()`, starting on line 206 of Listing 6.79. This subroutine is used by the actions for all of the logical operators. It is passed the numeric component of the target labels and a pointer to the operand's value structure. This second argument is used only to do a simple optimization—if the operand is already a temporary variable of type `int`, then it is recycled to hold the result of the operation, otherwise a new temporary variable is created. (The test is made on line 220 and the temporary is created on the next two lines.) Remember, both the test and the jump to the true or false label are generated before `gen_false_true()` is called, and once the jump has been created, the associated value is no longer needed. The operand's value could be discarded by the calling routine if you didn't want to bother with the optimization.

`gen_false_true()`

Note that the low-level `tmp_create()` function, which creates but doesn't initialize the temporary, is used here. Also note that the value is discarded before the `tmp_create()` call so that the same region of memory can be recycled for the result of the operation.

Listing 6.78. c.y— Unary Operators (Part 5)—Unary Arithmetic Operators

```
658 /* unop: */
659     | MINUS    unary          { $$ = do_unop( '-', $2 ); }      %prec UNOP
660     | UNOP     unary          { $$ = do_unop( $1, $2 ); }
```

The pre and post, increment and decrement operators are processed by the productions in Listing 6.80 and the `incop()` subroutine in Listing 6.81. The following input file:

`++ and --, incop().`

```
int x;
foo() { ++x; }
```

generates the following code for the `++x`:

```
W(&_x)    += 1;           /* Increment x itself. */
W( T(1) ) = W(&_x);      /* Copy incremented value to temporary. */
```

The synthesized attribute is the temporary variable. A postincrement like `x++` works in a similar manner, but the two instructions are transposed:

```
W( T(1) ) = W(&_x);
W(&_x)    += 1;
```

Again, the synthesized attribute is the temporary—the value of the temporary is determined by whether a pre or post increment is used. The `++` and `--` operators require an lvalue operand, but they both generate rvalue temporaries. Pre and postdecrements use `=-` rather than `+=` but are otherwise the same as the increments.

Listing 6.82 contains various right-hand sides for the address-of operator (`&`) and the dereference operators (`*`, `[]`, `.`, and `->`). The structure operators are grouped with the rest of them because they all do essentially the same thing: add an offset to a pointer and load the object at the indicated address into a temporary variable. All of these operators generate rvalues. The all require either an lvalue operand, or an rvalue operand of type “array”, “structure”, or “pointer”. Remember that aggregate-type references generate a temporary variable (an rvalue) holding the address of the object.

`Pointer and array dereferencing.`

The first issue is getting an objects address. The address-of operator is processed by `unop→AND unary` on line 665 of Listing 6.82. The `addr_of()` workhorse function is in Listing 6.83, below. Only two kinds of incoming values are legal here: an lvalue or an rvalue for an aggregate type such as a structure. In either case, the incoming value

`Address-of operator (&)`

Listing 6.79. *op.c*—Unary Minus and Ones-Complement Operator Processing

```

177 value *do_unop( op, val )
178 int op;
179 value *val;
180 {
181     char *op_buf = "=?";
182     int i;
183
184     if( op != '!' ) /* ~ or unary - */
185     {
186         op_buf[1] = op;
187         gen( op_buf, val->name, val->name );
188     }
189     else /* ! */
190     {
191         gen( "EQ", rvalue(val), "0" ); /* EQ(x, 0) */
192         gen( "goto%s%d", L_TRUE, i = tf_label() ); /* goto T000; */
193         val = gen_false_true( i, val ); /* fall thru to F */
194     }
195     return val;
196 }
197 /*-----*/
198
199 int tf_label()
200 {
201     static int label; /* Return the numeric component of a label for use as */
202     /* a true/false target for one of the statements */
203     return ++label; /* processed by gen_false_true(). */
204 }
205 /*-----*/
206
207 value *gen_false_true( labelnum, val )
208 int labelnum;
209 value *val;
210 {
211     /* Generate code to assign true or false to the temporary represented by
212      * val. Also, create a temporary to hold the result, using val if it's
213      * already the correct type. Return a pointer to the target. The FALSE
214      * code is at the top. If val is NULL, create a temporary. Labelnum must
215      * have been returned from a previous tf_label() call.
216 */
217     if( !val )
218         val = tmp_create( NULL, 0 );
219
220     else if( !val->is_tmp || !IS_INT(val->type) )
221     {
222         release_value( val );
223         val = tmp_create( NULL, 0 );
224     }
225     gen( ":%s%d", L_FALSE, labelnum ); /* F000: */ /* t0 = 0 */
226     gen( "=", val->name, "0" ); /* t0 = 0 */
227     gen( "goto%s%d", L_END, labelnum ); /* goto E000; */
228     gen( ":%s%d", L_TRUE, labelnum ); /* T000: */ /* t0 = 1; */
229     gen( "=", val->name, "1" ); /* t0 = 1; */
230     gen( ":%s%d", L_END, labelnum ); /* E000: */
231     return val;
232 }

```

Listing 6.80. *c.y*— Unary Operators (Part 6)—Pre and Postincrement

```

661 /* unop: */
662     | unary INCOP          { $$ = incop( 0, $2, $1 );      } %prec UNOP
663     | INCOP unary          { $$ = incop( 1, $1, $2 );      } %prec UNOP

```

Listing 6.81. *op.c*— Increment and Decrement Operators

```

233 value *incop( is_preincrement, op, val )           /* ++ or -- */
234 int   is_preincrement;                          /* Is preincrement or predecrement. */
235 int   op;                                     /* '-' for decrement, '+' for increment */
236 value *val;                                  /* lvalue to modify. */
237 {
238     char buf[ VALNAME_MAX ];
239     char *name;
240     value *new;
241     char *out_op = (op == '+') ? "+=%s%d" : "-=%s%d" ;
242     int inc_amt;
243
244     /* You must use rvalue_name() in the following code because rvalue()
245      * modifies the val itself--the name field might change. Here, you must use
246      * the same name both to create the temporary and do the increment so you
247      * can't let the name be modified.
248     */
249
250     if( !val->lvalue )
251         yyerror("%c%c: lvalue required\n", op, op );
252     else
253     {
254         inc_amt = (IS_POINTER(val->type)) ? get_sizeof(val->type->next) : 1 ;
255         name    = rvalue_name( val );
256
257         if( is_preincrement )
258         {
259             gen( out_op, name, inc_amt );
260             val = tmp_gen( val->type, val );
261         }
262         else                                /* Postincrement. */
263         {
264             val = tmp_gen( val->type, val );
265             gen( out_op, name, inc_amt );
266         }
267     }
268     return val;
269 }

```

Listing 6.82. *c.y*— Unary Operators (Part 7)—Indirection

```

664 /* unop: */
665     | AND      unary          { $$ = addr_of( $2 );      } %prec UNOP
666     | STAR     unary          { $$ = indirect( NULL, $2 ); } %prec UNOP
667     | unary   LB expr RB     { $$ = indirect( $3, $1 );      } %prec UNOP
668     | unary   STRUCTOP NAME { $$ = do_struct($1, $2, yytext); } %prec STRUCTOP

```

* and []

already holds the desired address: The lvalue is an address by definition. An rvalue like the one created by *unop*→NAME for an aggregate object is a physical temporary variable that also holds the required address. So, all that the address-of action needs to do is modify the value's type chain by adding an explicit pointer declarator at the far left and change its value into an rvalue by clearing the lvalue bit. Keep this action in mind when you read about the * operator, below. I'll start with the pointer dereference (*) and the array dereference ([]), operators, handled by indirect() (in Listing 6.84, below).

Listing 6.83. op.c—Address-of Operator Processing

```

270 value *addr_of( val )
271 value *val;
272 {
273     /* Process the & operator. Since the incoming value already holds the
274      * desired address, all you need do is change the type (add an explicit
275      * pointer at the far left) and change it into an rvalue. The first argument
276      * is returned.
277 */
278
279     link *p;
280
281     if( val->lvalue )
282     {
283         p           = new_link();
284         p->DCL_TYPE = POINTER;
285         p->next    = val->type;
286         val->type   = p;
287         val->lvalue = 0;
288     }
289     else if( !IS_AGGREGATE(val->type) )
290         yyerror( "(&) lvalue required\n" );
291
292     return val;
293 }
```

Listing 6.84. op.c—Array Access and Pointer Dereferencing

```

294 value *indirect( offset, ptr )          /* Offset factor (NULL if it's a pointer). */
295 value *offset;                         /* Pointer that holds base address. */
296 value *ptr;                           */
297 {
298     /* Do the indirection, If no offset is given, we're doing a *, otherwise
299      * we're doing ptr[offset]. Note that, strictly speaking, you should create
300      * the following dumb code:
301      *
302      *          t0 = rvalue( ptr );      (if ptr isn't a temporary)
303      *          t0 += offset        (if doing [offset])
304      *          t1 = *t0;           (creates a rvalue)
305      *                               lvalue attribute = &t1
306      *
307      * but the first instruction is necessary only if ptr is not already a
308      * temporary, the second only if we're processing square brackets.
309      *
310      * The last two operations cancel if the input is a pointer to a pointer. In
311      * this case all you need to do is remove one * declarator from the type
312      * chain. Otherwise, you have to create a temporary to take care of the type
313      * conversion.
314 */
```

Listing 6.84. continued...

```

315
316     link *tmp      ;
317     value *synth   ;
318     int  objsize   ; /* Size of object pointed to (in bytes)           */
319
320     if( !IS_PTR_TYPE(ptr->type) )
321         yyerror( "Operand for * or [N] Must be pointer type\n" );
322
323     rvalue( ptr );    /* Convert to rvalue internally. The "name" field   */
324     /* is modified by removing leading &'s from logical */          */
325     /* lvalues.                                              */
326     if( offset )      /* Process an array offset.                         */
327     {
328         if( !IS_INT(offset->type) || !IS_CHAR(offset->type) )
329             yyerror( "Array index must be an integral type\n" );
330
331         objsize = get_sizeof( ptr->type->next );      /* Size of dereferenced */
332                                         /* object.                */
333
334         if( !ptr->is_tmp )                                /* Generate a physical */
335             ptr = tmp_gen( ptr->type, ptr );            /* lvalue.               */
336
337         if( IS_CONSTANT( offset->type ) )                 /* Offset is a constant.*/
338         {
339             gen("+=%s%d", ptr->name, offset->type->V_INT * objsize );
340         }
341     else
342     {
343         /* Offset is not a constant. Do the arithmetic at run time. */
344
345         if( objsize != 1 )                                /* Multiply offset by */
346             /* size of one object. */
347         if( !offset->is_tmp )
348             offset = tmp_gen( offset->type, offset );
349
350             gen( "*=%s%d", offset->name, objsize );
351         }
352
353         gen( "+=",  ptr->name, offset->name );        /* Add offset to base. */
354     }
355
356     release_value( offset );
357 }
358
359 /* The temporary just generated (or the input variable if no temporary
360 * was generated) now holds the address of the desired cell. This command
361 * must generate an lvalue unless the object pointed to is an aggregate,
362 * whereupon it's an rvalue. In any event, you can just label the current
363 * cell as an lvalue or rvalue as appropriate and continue. The type must
364 * be advanced one notch to compensate for the indirection, however.
365 */
366
367     synth      = ptr;
368     tmp       = ptr->type;                            /* Advance type one notch. */
369     ptr->type = ptr->type->next;
370     discard_link(tmp);
371

```



Listing 6.84. continued...

```

372     if( !IS_AGGREGATE(ptr->type->next) )
373         synth->lvalue = 1;                         /* Convert to lvalue. */
374
375     return synth;
376 }
```

Operand to * or []
must be array or pointer.

Attribute synthesized by
* and [] operators.

Rules for forming lvalues
and rvalues when pro-
cessing * and [].

The best way to understand the code is to analyze what actually happens as the operator is processed. First of all, as was the case with the address-of operator, the operand must represent an address. It is one of two things: (1) **an array**, in which case the operand is an rvalue that holds the base address of the array, or (2) **a pointer**, in which case the operand is an lvalue for something of pointer type—the expression in the value's name field evaluates to the address of the pointer.

The next issue is the synthesized attribute, which is controlled by the type of the dereferenced object. The compiler needs to convert the incoming attribute to the synthesized attribute. If the dereferenced object is a nonaggregate, the synthesized attribute is an lvalue that holds the address of the dereferenced object. If the dereferenced object is an aggregate, then the attribute is an rvalue that holds the address of the first element of the object. In other words, the generated attribute must be the same thing that would be created by *unop*→NAME if it were given an identifier of the same type as the referenced object. Note that both the incoming and outgoing attributes are addresses.

To summarize (if the following isn't clear, hold on for a second—several examples follow):

- (1) The dereferenced object is not an aggregate. The synthesized attribute is an lvalue that references that object, and:
 - a. if the incoming attribute is an rvalue, then it already contains the necessary address. That is, the incoming rvalue holds the address of the dereferenced object, and the outgoing lvalue must also hold the address of the dereferenced object. Consequently, no code needs to be generated in this case. The compiler does have to modify the value internally, however. First, by setting the lvalue bit true, and second, by removing the first link in the type chain. Remember, the generated lvalue is for the dereferenced object. If the incoming type is “pointer to **int**”, the outgoing object is an lvalue that references the **int**.
 - If the compiler's doing an array access rather than a pointer access, code must also be emitted to add an offset to the base address that's stored in the lvalue.
 - b. if the incoming attribute is a logical lvalue, then all the compiler needs to do is remove the leading ampersand and change the type as discussed earlier. The pointer variable is now treated as if it were a physical lvalue.
 - c. if the incoming attribute is a physical lvalue, for a pointer, then you need to generate code to get rid of one level of indirection, and change the type as was discussed earlier. You can safely add an offset to the physical lvalue to do array access.
- (2) The dereferenced object is an aggregate. The synthesized attribute is an rvalue that points at the first element of the referenced array or structure, and:
 - a. if the incoming attribute is an rvalue, use it for the synthesized attribute, changing the type by removing the first declarator link and adding an offset as necessary.

- b. if the incoming attribute is a logical lvalue, create a physical rvalue, with the type adjusted, as above.
- c. if the incoming attribute is a physical lvalue, convert it to an rvalue by resetting the lvalue bit and adjust the type, as above.

Note that many of the foregoing operations generate no code at all—they just modify the way that something is represented internally. Code is generated only when an offset needs to be computed or an incoming lvalue references an aggregate object. Also, notice the similarities between processing a `&` and processing a `*`. The former adds a pointer declarator to the type chain and turns the incoming attribute into an rvalue; the latter does the exact opposite, removing a pointer declarator from the front of the type chain and turning the incoming attribute into an lvalue.

I'll demonstrate how indirection works with some examples. Parses of the following code is shown in Tables 6.17 to 6.23.

```
int *p, a[10];
foo()
{
    *p++;      *++p;      ++*p;      (*p)++;
    a[7];      ++p[3];   p++[3];
}
```

Table 6.17. Generate Code for `*p++;`

Parse and Value Stacks	Comments
<code>stmt_list</code> □	Shift STAR .
<code>stmt_list</code> STAR □ □	Shift NAME . The attribute for NAME is a pointer to the symbol for <code>p</code> .
<code>stmt_list</code> STAR NAME □ □ <code>p</code>	Reduce by unary → NAME . The action creates an lvalue that references the variable.
<code>stmt_list</code> STAR unary □ □ <code>&WP(&_p)_L</code>	Shift INCOP .
<code>stmt_list</code> STAR unary INCOP □ □ <code>&WP(&_p)_L</code> <code>'+'</code>	Reduce by unary → unary INCOP . Emit code to increment <code>p</code> : CODE: <code>WP(T(1)) = WP(&p);</code> CODE: <code>WP(&p) += 2;</code> The synthesized attribute is an rvalue (of type “pointer to <code>int</code> ”) for the temporary variable.
<code>stmt_list</code> STAR unary □ □ <code>WP(T(1))_R</code>	Reduce by unary → STAR unary . <code>\$2</code> (under the unary) already contains the correct address and is of the correct type. Convert it to a physical lvalue that references <code>*p</code> (not <code>p</code> itself).
<code>stmt_list</code> unary □ <code>WP(T(1))_L</code>	Subsequent operations cannot access <code>p</code> ; they can access <code>*p</code> with <code>*WP(T(1))</code> .

Neither the entire parse nor the entire parse stack is shown in the tables, but there's enough shown that you can see what's going on. The following symbols are on the stack to the left of the ones that are shown.

`ext_def_list opt_specifiers funct_decl {70} def_list {71} LC {65} local_defs`

I've shown the value stacks under the parse stack. Attributes for most of the nonterminals are value structures—the name field is shown in the table. Subscripts indicate whether it's an lvalue (`name`) or rvalue (`name`). Logical lvalues start with an ampersand. A box (`□`) is used when there are no attributes. You should look at all these tables now, even though they're scattered over several pages.

I've used the `++` operators in these expressions because `++` requires an lvalue, but it generates an rvalue. The `++` is used to demonstrate how the indirection is handled with both kinds of incoming attributes. Other, more complex expressions [like `*(p+1)`] are

Table 6.18. Generate Code for $\star\star p;$

Parse and Value Stacks	Comments
<i>stmt_list</i> □	Shift STAR
<i>stmt_list</i> □ STAR □	Shift INCOP . The attribute for INCOP is the first character of the lexeme.
<i>stmt_list</i> □ STAR INCOP □ □ '+'	Shift NAME . The attribute for NAME is a pointer to the symbol for <i>p</i> .
<i>stmt_list</i> □ STAR INCOP NAME □ □ '+' <i>p</i>	Reduce by <i>unary</i> → NAME . Create a logical lvalue that references <i>p</i> .
<i>stmt_list</i> □ STAR INCOP <i>unary</i> □ □ '+' &WP (& <i>p</i>) _L	Reduce by <i>unary</i> → INCOP unary . Emit code to increment <i>p</i> : CODE: WP (& <i>p</i>) += 2; CODE: WP (T(1)) = WP (& <i>p</i>); The synthesized attribute is an rvalue of type “pointer to int ” that holds a copy of <i>p</i> . From this point on, the compiler has forgotten that <i>p</i> ever existed—at least for the purposes of processing the current expression.
<i>stmt_list</i> □ STAR <i>unary</i> □ WP (T(1)) _R	Reduce by <i>unary</i> → STAR unary . The current reduction turns the attribute from the previous reduction it into a physical lvalue that references the dereferenced object (<i>*p</i>).
<i>stmt_list</i> □ <i>unary</i> WP (T(1)) _L	

Table 6.19. Generate Code for $(\star p)++;$

Parse and Value Stacks	Comments
<i>stmt_list</i> □	Shift LP
<i>stmt_list</i> □ LP □	Shift STAR
<i>stmt_list</i> □ LP STAR □ □	Shift NAME . The attribute is a pointer to a symbol structure for <i>p</i> .
<i>stmt_list</i> □ LP STAR NAME □ □ □ <i>p</i>	Reduce by <i>unary</i> → NAME . The synthesized attribute is a logical lvalue that references <i>p</i> .
<i>stmt_list</i> □ LP STAR <i>unary</i> □ □ □ &WP (& <i>p</i>) _L	Reduce by <i>unary</i> → STAR unary . Convert the incoming logical lvalue for <i>p</i> to a physical lvalue that references <i>*p</i> .
<i>stmt_list</i> □ LP <i>unary</i> □ WP (& <i>p</i>) _L	<i>unary</i> is now converted to <i>expr</i> by a series of reductions, not shown here. The initial attribute is passed through to the <i>expr</i> , however.
<i>stmt_list</i> □ LP <i>expr</i> □ WP (& <i>p</i>) _L	Shift RP
<i>stmt_list</i> □ LP <i>expr</i> RP □ □ WP (& <i>p</i>) _L	Reduce by <i>unary</i> → LP expr RP . \$\$ = \$2;
<i>stmt_list</i> □ <i>unary</i> WP (& <i>p</i>) _L	Shift INCOP
<i>stmt_list</i> □ <i>unary</i> INCOP WP (& <i>p</i>) _L '+'	Reduce by <i>unary</i> → <i>unaryINCOP</i> . Emit code to increment <i>*p</i> : CODE: W(T(1)) = *WP (& <i>p</i>); CODE: *WP (& <i>p</i>) += 1; The synthesized attribute is an rvalue for the temporary that holds the incremented value of <i>*p</i> ;
<i>stmt_list</i> □ <i>unary</i> W(T(1)) _R	

handled in much the same way. (*p*+1, like $\star\star p$, generates an rvalue of type pointer.)

I suggest doing an exercise at this juncture. Parse the following code by hand, starting at the *stmt_list*, as in the previous examples, showing the generated code and the relevant parts of the parse and value stack at each step:

Table 6.20. Generate Code for $++^*p;$

Parse and Value Stacks	Comments
<i>stmt_list</i> □	Shift INCOP
<i>stmt_list</i> □ INCOP □ '+'	Shift STAR
<i>stmt_list</i> □ INCOP STAR □ '+' □	Shift NAME. The attribute for NAME is a pointer to the symbol structure for p.
<i>stmt_list</i> □ INCOP STAR NAME □ '+' □ p	Reduce by $unary \rightarrow NAME$. Create a logical lvalue for p.
<i>stmt_list</i> □ INCOP STAR unary □ '+' □ &WP(&_p)_L	Reduce by $unary \rightarrow STAR\ unary$. This reduction converts the logical lvalue that references p into a physical lvalue that references *p. CODE: *WP(&_p) += 1; CODE: W(T(1)) = *WP(&_p);
<i>stmt_list</i> □ INCOP unary □ '+' □ WP(&_p)_L	Reduce by $unary \rightarrow INCOP\ unary$ Emit code to increment *p: CODE: *WP(&_p) += 1; The generated attribute is an rvalue that holds the contents of *p after the increment—this is a preincrement.
<i>stmt_list</i> □ unary □ W(T(1))_R	

Table 6.21. Generate Code for $a[7];$

Parse and Value Stacks	Comments
<i>stmt_list</i> □	Shift NAME. The attribute for NAME is a pointer to the symbol structure for a.
<i>stmt_list</i> □ NAME □ a	Reduce by $unary \rightarrow NAME$. Since this is an array rather than a simple pointer variable, an rvalue is generated of type “pointer to first array element” (pointer to int , here): CODE: W(T(1)) = &W(_a);
<i>stmt_list</i> □ unary □ W(T(1))_R	Shift LB.
<i>stmt_list</i> □ unary LB □ W(T(1))_R □	Shift ICON.
<i>stmt_list</i> □ unary LB ICON □ W(T(1))_R □ "7"	Reduce by $unary \rightarrow ICON$. The synthesized attribute is a value representing the constant 7.
<i>stmt_list</i> □ unary LB unary □ W(T(1))_R □ 7_R	$unary$ is reduced to $expr$ by a series of reductions, not shown here. The attribute is passed through to the $expr$, however.
<i>stmt_list</i> □ unary LB expr □ W(T(1))_R □ 7_R	Shift RB.
<i>stmt_list</i> □ unary LB expr RB □ W(T(1))_R □ 7_R □	Reduce by $unary \rightarrow LBexprRB$. The incoming attribute is an rvalue of type “pointer to int .” The following code is generated to compute the offset: CODE: W(T(1)) += 14; The synthesized attribute is an lvalue holding the address of the eighth array element.
<i>stmt_list</i> □ unary □ W(T(1))_L	The final attribute is an lvalue that references a[7].

```
int *x[5];
**x;
*x[3];
x[1][2];
```

The other pointer-related operators are the structure-access operators: `.` and `->`. These are handled by the `do_struct()` subroutine in Listing 6.85. The basic logic is the same as that used for arrays and pointers. This time the offset is determined by the position of the field within the structure, however.

The final unary operator handles function calls. It's implemented by the `unop` right-hand sides in Listing 6.85 and the associated subroutines, `call()` and `ret_reg()`, in Listing 6.86. One of the more interesting design issues in C is that the argument list as a

Structure access: `->` and `.`, `do_struct()`.

Function calls.

Table 6.22. Generate Code for `p++[3];`

Parse and Value Stacks	Comments
<i>stmt_list</i> □	Shift NAME
<i>stmt_list</i> <i>NAME</i> □	Reduce by <i>unary</i> → NAME .
<i>stmt_list</i> <i>unary</i> □ &WP (&_p) _L	Shift INCOP
<i>stmt_list</i> <i>unary</i> INCOP □ &WP (&_p) _L '+'	Reduce by <i>unary</i> → <i>unary</i> INCOP . Emit code to increment the pointer. The synthesized attribute is an rvalue of type pointer to int . CODE: WP(T(1)) = WP(&_p); CODE: WP(&_p) += 2;
<i>stmt_list</i> <i>unary</i> □ WP(T(1)) _R	Shift LB .
<i>stmt_list</i> <i>unary</i> LB □ WP(T(1)) _R □	Shift ICON .
<i>stmt_list</i> <i>unary</i> LB ICON □ WP(T(1)) _R □ '3'	Reduce by <i>unary</i> → ICON . The action creates an rvalue for an int constant. The numeric value is stored in <code>\$\$->etype->V_INT</code> .
<i>stmt_list</i> <i>unary</i> LB <i>unary</i> □ WP(T(1)) _R □ 3 _R	Reduce <i>unary</i> to <i>expr</i> by a series of reductions, not shown. The original attribute is passed through to the <i>expr</i> .
<i>stmt_list</i> <i>unary</i> LB <i>expr</i> □ WP(T(1)) _R □ 3 _R	Shift RB .
<i>stmt_list</i> <i>unary</i> LB <i>expr</i> RB □ WP(T(1)) _L □ 3 _R □	Reduce by <i>unary</i> → ICON . Code is generated to compute the address of the fourth cell: CODE: WP(T(1)) += 6; Since the incoming attribute is already a temporary variable of the correct type, there's no need to generate a second temporary here—the first one is recycled. Note that it is translated to an lvalue that references the fourth cell, however.
<i>stmt_list</i> <i>unary</i> □ WP(T(1)) _L	

whole can be seen as a function-call operator which can be applied to any function pointer,²² a function name evaluates to a pointer to a function, much like an array name evaluates to a pointer to the first array element. The code-generation action does four things: push the arguments in reverse order, call the function, copy the return value into an rvalue, discard the arguments by adding a constant to the stack pointer. For example, a call like the following:

```
doctor( lawyer, merchant, chief );
```

is translated to:

```
push( W(&_chief) ); /* push the arguments */  

push( W(&_merchant) );  

push( W(&_lawyer) );  

call( _doctor ); /* call the subroutine */  

W( T(1) ) = rF.w.low; /* copy return value to rvalue */  

sp += 3; /* discard arguments. */
```

Function-argument processing,
non_comma_expr.

The function arguments are processed by the *args* productions on lines 674 to 681 of Listing 6.86. A *non_comma_expr* recognizes all C expressions except those that use comma operators. You must do something like the following to get a comma operator into a function call:

22. This is actually a C++ism, but it's a handy way to look at it.

Table 6.23. Generate Code for `++p[3];`

Parse and Value Stacks				Comments
<i>stmt_list</i> □				Shift INCOP
<i>stmt_list</i> □ INCOP '+'				Shift NAME. The attribute is a symbol representing p.
<i>stmt_list</i> □ INCOP NAME '+'				Reduce by <code>unary→NAME</code> . Convert symbol to an lvalue.
<i>stmt_list</i> □ INCOP unary '+' &WP(&_p),				Shift LB
<i>stmt_list</i> □ INCOP unary LB '+' &WP(&_p), □				Shift ICON
<i>stmt_list</i> □ INCOP unary LB ICON '+' &WP(&_p), □ '3'				Reduce by <code>unary→ICON</code> . The synthesized attribute is a value structure representing an <code>int</code> constant. It's an rvalue. The actual value (3) is stored internally in that structure.
<i>stmt_list</i> □ INCOP unary LB unary '+' &WP(&_p), □ 3 _R				Reduce <code>unary</code> to <code>expr</code> by a series of reductions, not shown. The original attribute is passed through to the <code>expr</code> .
<i>stmt_list</i> □ INCOP unary LB expr '+' &WP(&_p), □ 3 _R				Shift RB.
<i>stmt_list</i> □ INCOP unary LB expr RB '+' &WP(&_p), □ 3 _R □				Reduce by <code>unary→unaryLB expr RB</code> . Code is generated to compute the address of the fourth cell, using the offset in the value at \$3 and the base address in the value at \$1: <code>CODE: WP(T(1)) = WP(&_p);</code> <code>CODE: WP(T(1)) += 6;</code> You have to generate a physical lvalue because p itself may not be incremented. The synthesized attribute is an lvalue that holds the address of the fourth element.
<i>stmt_list</i> □ INCOP unary '+' WP(T(1)),				Reduce by <code>unary→INCOP unary</code> . Code is generated to increment the array element, the address of which is in the lvalue generated in the previous reduction. <code>CODE: *WP(T(1)) += 1;</code> <code>CODE: W(T(2)) = *WP(T(1));</code> The synthesized attribute is an rvalue that duplicates the contents of that cell.
<i>stmt_list</i> □ unary W(T(2)) _R				

```
tinker( (tailor,cowboy), sailor );
```

This function call has two arguments, the first one is the expression `(tailor, cowboy)`, which evaluates to `cowboy`. The second argument is `sailor`. The associated attribute is a value for that expression—an lvalue is used if the expression is a simple variable or pointer reference; an rvalue is used for most other expressions.

The `args` productions just traverse the list of arguments, printing the push instructions and keeping track of the number of arguments pushed. The argument count is returned back up the tree as the synthesized attribute. Note that the `args` productions form a right-recursive list. Right recursion is generally not a great idea in a bottom-up parser because all of the list elements pile up on the stack before any reductions occur. On the other hand, the list elements are processed in right to left order, which is convenient here because arguments have to be pushed from right to left. The recursion shouldn't cause problems unless the subroutine has an abnormally high number of arguments.

The `call()` subroutine at the top of Listing 6.87 generates both the `call` instruction and the code that handles return values and stack clean up. It also takes care of implicit subroutine declarations on lines 513 to 526. The action in `unary→NAME` creates a symbol of type `int` for an undeclared identifier, and this symbol eventually ends up here as

`args→...` productions.
Right recursion gets arguments pushed in correct order.

`call()`
`unary→NAME`

the incoming attribute. The `call()` subroutine changes the type to “function returning int” by adding another link to the head of the type chain. It also clears the `implicit` bit to indicate that the symbol is a legal implicit declaration rather than an undeclared variable. Finally, a C-code `extern` statement is generated for the function.

Listing 6.85. *op.c*—Structure Access

```

377 value *do_struct( val, op, field_name )
378 value *val;
379 int    op;           /* . or - (the last is for ->) */
380 char   *field_name;
381 {
382     value      *new;
383     symbol     *field;
384     link       *lp;
385
386     /* Structure names generate rvalues of type structure. The associated      */
387     /* name evaluates to the structure's address, however. Pointers generate      */
388     /* lvalues, but are otherwise the same.                                         */
389
390     if( IS_POINTER(val->type) )
391     {
392         if( op != '-' )
393         {
394             yyerror("Object to left of -> must be a pointer\n");
395             return val;
396         }
397         lp      = val->type;          /* Remove POINTER declarator from */
398         val->type = val->type->next;  /* the type chain and discard it. */
399         discard_link( lp );
400     }
401
402     if( !IS_STRUCT(val->type) )
403     {
404         yyerror("Not a structure.\n");
405         return val;
406     }                                /* Look up the field in the structure table: */
407
408     if( !(field = find_field(val->type->V_STRUCT, field_name)) )
409     {
410         yyerror("%s not a field\n", field_name );
411         return val;
412     }
413
414     if( val->lvalue || !val->is_tmp )          /* Generate temporary for      */
415         val = tmp_gen( val->type, val );        /* base address if necessary; */
416                                              /* then add the offset to the */
417     if( field->level > 0 )                      /* desired field.               */
418         gen( "+=%s%d", val->name, field->level );
419
420     if( !IS_AGGREGATE(field->type) )           /* If referenced object isn't */
421         val->lvalue = 1;                         /* an aggregate, use lvalue. */
422
423                                              /* Replace value's type chain */
424                                              /* with type chain for the */
425                                              /* referenced object:        */
426     lp = val->type;

```



Listing 6.85. continued...

```

427     if( !(val->type = clone_type( field->type, &val->etype)) )
428     {
429         yyerror("INTERNAL do_struct: Bad type chain\n" );
430         exit(1);
431     }
432     discard_link_chain( lp );
433     access_with( val );                                /* Change the value's name */
434                                         /* field to access an object */
435                                         /* of the new type. */          */
436 }
437
438 /*-----*/
439
440 PRIVATE symbol *find_field( s, field_name )
441 structdef      *s;
442 char           *field_name;
443 {
444     /* Search for "field_name" in the linked list of fields for the input
445      * structdef. Return a pointer to the associated "symbol" if the field
446      * is there, otherwise return NULL.
447      */
448
449     symbol      *sym;
450     for( sym = s->fields; sym; sym = sym->next )
451     {
452         if( !strcmp(field_name, sym->name) )
453             return sym;
454     }
455     return NULL;
456 }
457
458 /*-----*/
459
460 PRIVATE char    *access_with( val )
461 value      *val;
462 {
463     /* Modifies the name string in val so that it references the current type.
464      * Returns a pointer to the modified string. Only the type part of the
465      * name is changed. For example, if the input name is "WP(fp+4)", and the
466      * type chain is for an int, the name is be changed to "W(fp+4). If val is
467      * an lvalue, prefix an ampersand to the name as well.
468      */
469
470     char *p, buf[ VALNAME_MAX ] ;
471
472     strcpy( buf, val->name );
473     for( p = buf; *p && *p != '(' /* */ ; ++p )           /* find name */
474     ;
475
476     if( !*p )
477         yyerror( "INTERNAL, access_with: missing parenthesis\n" );
478     else
479         sprintf( val->name, "%s%s%s", val->lvalue ? "&" : "",      get_prefix(val->type), p );
480
481     return val->name;
482 }
```

Listing 6.86. *c.y*— Unary Operators (Part 8)—Function Calls

```

669 /* unop: */
670     | unary LP args RP      { $$ = call      ( $1,   $3 );    }
671     | unary LP           RP      { $$ = call      ( $1,   0  );    }
672     ;
673
674 args     : non_comma_expr %prec COMMA     { gen( "push", rvalue( $1 ) );
675                                         release_value( $1 );
676                                         $$ = 1;
677                                         }
678     | non_comma_expr COMMA args     { gen( "push", rvalue( $1 ) );
679                                         release_value( $1 );
680                                         $$ = $3 + 1;
681                                         }
682     ;

```

Listing 6.87. *op.c*— Function-Call Processing

```

483 PUBLIC value *call( val, nargs )
484 value *val;
485 int   nargs;
486 {
487     link  *lp;
488     value *synth;          /* synthesized attribute */ 
489
490     /* The incoming attribute is an lvalue for a function if
491      *             funct()
492      * or
493      *             int (*p) () = funct;
494      *             (*p) ();
495      *
496      * is processed. It's a pointer to a function if p() is used directly.
497      * In the case of a logical lvalue (with a leading &), the name will be a
498      * function name, and the rvalue can be generated in the normal way by
499      * removing the &. In the case of a physical lvalue the name of a variable
500      * that holds the function's address is given. No star may be added.
501      * If val is an rvalue, then it will never have a leading &.
502      */
503
504     if( val->sym  &&  val->sym->implicit  && !IS_FUNCT(val->type) )
505     {
506         /* Implicit symbols are not declared. This must be an implicit function
507          * declaration, so pretend that it's explicit. You have to modify both
508          * the value structure and the original symbol because the type in the
509          * value structure is a copy of the original. Once the modification is
510          * made, the implicit bit can be turned off.
511          */
512
513     lp                  = new_link();
514     lp->DCL_TYPE       = FUNCTION;
515     lp->next            = val->type;
516     val->type           = lp;
517
518     lp                  = new_link();
519     lp->DCL_TYPE       = FUNCTION;
520     lp->next            = val->sym->type;
521     val->sym->type     = lp;
522

```

Listing 6.87. continued...

```

523         val->sym->implicit = 0;
524         val->sym->level    = 0;
525
526         yydata( "extern\t%s();\n", val->sym->rname );
527     }
528
529     if( !IS_FUNCT(val->type) )
530     {
531         yyerror( "%s not a function\n", val->name );
532         synth = val;
533     }
534     else
535     {
536         lp      = val->type->next;           /* return-value type */
537         synth = tmp_create( lp, 0 );
538
539         gen( "call", *val->name == '&' ? &val->name[1] : val->name );
540         gen( "=", synth->name, ret_reg(lp) );
541
542         if( nargs )
543             gen( "+=%s%d" , "sp", nargs );
544
545         release_value( val );
546     }
547
548     return synth;
549 }
550
551 -----*/
552
553 char   *ret_reg( p )
554 link   *p;
555 {
556     /* Return a string representing the register used for a return value of
557      * the given type.
558      */
559
560     if( IS_DECLARATOR( p ) )
561         return "rF.pp";
562     else
563         switch( p->NOUN )
564         {
565             case INT:      return (p->LONG) ? "rF.l" : "rF.w.low" ;
566             case CHAR:    return "rF.w.low" ;
567             default:    yyerror("INTERNAL ERROR: ret_reg, bad noun\n");
568                         return "AAAAAAAAAAAAAGH!";
569         }
570     }

```

6.8.5 Binary Operators

We can now breathe a collective sigh of relief. The declaration system and the unary operators are the hardest parts of the compiler and they're over with. We can now move on to the binary operators. Again, it's useful to look at the overall structure of the binary-operator productions before proceeding. These are summarized in Table 6.24.

Table 6.24. Summary of Binary-Operator Productions

<i>expr</i>	:	<i>expr</i>	COMMA	<i>non_comma_expr</i>		
		<i>non_comma_expr</i>				
<i>non_comma_expr</i>	→	<i>non_comma_expr</i>	ASSIGNOP	<i>non_comma_expr</i>		
		<i>non_comma_expr</i>	EQUAL	<i>non_comma_expr</i>		
		<i>non_comma_expr</i>	QUEST	<i>non_comma_expr</i>	COLON	<i>non_comma_expr</i>
		<i>or_expr</i>				
<i>or_expr</i>	→	<i>or_list</i>				
<i>or_list</i>	→	<i>or_list</i>	OROR	<i>and_expr</i>		
		<i>and_expr</i>				
<i>and_expr</i>	→	<i>and_list</i>				
<i>and_list</i>	→	<i>and_list</i>	ANDAND	<i>binary</i>		
		<i>binary</i>				
<i>binary</i>	→	<i>binary</i>	PLUS	<i>binary</i>		
		<i>binary</i>	MINUS	<i>binary</i>		
		<i>binary</i>	STAR	<i>binary</i>		
		<i>binary</i>	DIVOP	<i>binary</i>		
		<i>binary</i>	SHIFTOP	<i>binary</i>		
		<i>binary</i>	AND	<i>binary</i>		
		<i>binary</i>	XOR	<i>binary</i>		
		<i>binary</i>	OR	<i>binary</i>		
		<i>binary</i>	RELOP	<i>binary</i>		
		<i>binary</i>	EQUOP	<i>binary</i>		
		<i>unary</i>				

This grouping of the productions is a compromise between the strict grammatical approach in which the grammar alone determines precedence and associativity—all operators that share a precedence level are isolated into single productions—and the more flexible **yacc**/**occs** approach in which all binary operators are combined into a single production with many right-hand sides, associativity and precedence controlled by `%left` and `%right` directives. As you can see, the **yacc** approach has been used for most of the operators by putting them under the aegis of a single *binary* production. The `%left` and `%right` directives that control the precedence and associativity of these operators were discussed back towards the beginning of the current chapter. They're also at the top of Appendix C if you want to review them.

Several of the lower-precedence operators have been isolated from the rest of the binary operators for grammatical and code-generation reasons, however. The first issue is the *non_comma_expr* which recognizes all expressions except those that use the comma operator at the outermost parenthesis-nesting level. The following statement is not recognized by a *non_comma_expr*:

```
a = foo(), bar();
```

but this one is okay:

```
(a = foo(), bar());
```

because of the parenthesis. (Neither expression is a model of good programming style, but the language accepts them.) An *expr* accepts the comma operator at the outer level; both of the foregoing expressions are recognized. This isolation is required for grammatical reasons. Comma-delimited lists of expressions are used in several places (for subroutine-argument lists, initializer lists, and so on). If an unparenthesised comma operator were permitted in any of these comma-delimited lists, the parser would not be able to distinguish between a comma that separated list elements and the comma operator. The practical consequence would be reduce/reduce conflicts all over the place. Isolating the comma operator lets you use a *non_comma_expr* in all of the comma-

Isolating the comma operator. Problems with comma-delimited lists.

delimited lists and an *expr* everywhere else.

The *expr* production, which handles the comma operator, is at the top of Listing 6.88. The imbedded action releases the value structure associated with the left element of the expression; the value for the right element is used as the synthesized attribute. The imbedded action is necessary to prevent the compiler from wasting a temporary variable. If \$1 was released at the far right of the production, then both components of the list would be evaluated before the left one was discarded. (Run a parse of `++a, ++b` by hand if you don't believe me.) Consequently, the temporary that holds the evaluated result of the *expr* would continue to exist while the *non_comma_expr* was being evaluated. Since the left term is going to be discarded anyway, there's no point in keeping it around longer than necessary.

The comma operator.

Listing 6.88. *c.y*— Binary Operators: Comma, Conditional, and Assignment

```

683 expr      : expr COMMA {release_value($1);}  non_comma_expr { $$=$4; }
684     | non_comma_expr
685 ;
686
687 non_comma_expr
688     : non_comma_expr QUEST { static int label = 0;
689
690             gen( "EQ",      rvalue( $1 ), "0" );
691             gen( "goto%s%d", L_COND_FALSE,
692                               $<num>$ = ++label );
693             release_value( $1 );
694         }
695     non_comma_expr COLON { $<p_val>$ = $4->is_tmp
696                             ? $4
697                             : tmp_gen($4->type, $4);
698
699             gen( "goto%s%d", L_COND_END,   $<num>3 );
700             gen( ":%s%d",    L_COND_FALSE, $<num>3 );
701
702         }
703     non_comma_expr { $$ = $<p_val>6;
704
705         if( !the_same_type($$->type, $7->type, 1) )
706             yyerror("Types on two sides of colon \"\
707                      \"must agree\n");
708
709         gen( "=",      $$->name,      rvalue($7) );
710         gen( ":%s%d", L_COND_END,   $<num>3 );
711         release_value( $7 );
712     }
713
714     | non_comma_expr ASSIGNOP non_comma_expr { $$ = assignment($2, $1, $3); }
715     | non_comma_expr EQUAL      non_comma_expr { $$ = assignment( 0, $1, $3); }
716
717     | or_expr
718 ;

```

The *non_comma_expr* production, also in Listing 6.88, handles the conditional and assignment operators. A statement like this:

The conditional operator
(*a*?*b*:*c*).

```
int mickey, minnie;
mickey = minnie ? pluto() : goofy();
```

generates the code in Listing 6.89.

Listing 6.89. Code Generated for the Conditional Operator.

```

1      EQ( W(&_minnie), 0 )    /* if( minnie == 0 )                      */
2      goto QF1;             /*      branch around the first clause   */
3
4      call(_pluto);        /* true part of the conditional:       */
5      W( T(1) ) = rF.w.low; /*      rvalue = subroutine return value */
6      goto QE1;
7  QF1:
8      call(_goofy);        /* false part of the conditional */
9      W( T(2) ) = rF.w.low; /*      rvalue = subroutine return value */
10
11     W( T(1) ) = W( T(2) );
12  QE1:
13     W(&_mickey) = W( T(1) ); /* final assignment */

```

The conditional is processed by the right-hand side on lines 688 to 712 of Listing 6.88. The only tricky issue here is the extra assignment just above the QE1 label in Listing 6.89. The problem, here, is inherent in the way a bottom-up parser works. It's very difficult for a bottom-up parser to tell the expression-processing code to put the result in a particular place; rather, the expression-processing code decides more or less arbitrarily where the final result of an expression evaluation will be found, and it passes that information back up the parse tree to the higher-level code. The difficulty with the conditional operator is the two action clauses (one for the true condition and a second for the false condition), both of which generate a temporary variable holding the result of the expression evaluation. Since the entire conditional must evaluate to only one temporary, code must be generated to copy the value returned from the false clause into the same temporary that was used for the true clause.

None of this would be a problem with a top-down parser, which can tell the subroutine that processes the action clause to put the final result in a particular place. For example, the high-level subroutine that processes the conditional in a recursive-descent parser can pass a value structure to the lower-level subroutines that generate the expression-processing code. These lower-level routines could then use that value for the final result of the expression-evaluation process.

Assignment operators.

The other two right-hand sides to *non_comma_expr* (on lines 714 and 715 of Listing 6.89) handle the assignment operators. The *assignment()* function in Listing 6.90 does all the work. The subroutine is passed three arguments: an operator (*op*) and values for the destination (*dst*) and source (*src*). The operator argument is zero for simple assignment, otherwise it's the first character of the lexeme: '+' for *+=*, '<' for *<=*, and so on. This operator is, in turn, passed to *gen()* by adding it to the operator string on lines 588 and 589 of Listing 6.90.

Implicit type conversion.

The destination argument must be an lvalue. If it's a physical lvalue, *assignment()* tells *gen()* to add a star to the right of the destination name by adding an @ to the operator string on line 587 of Listing 6.90. The source operand is converted to an rvalue on line 593, then the assignment code is generated.

If the source and destination types don't match, an implicit type conversion must be performed as part of the assignment on line 616. The *convert_type()* subroutine is the same one that was used for temporary-variable initialization. (It's in Listing 6.69 on page 590.) *convert_type()* generates code to copy the source variable into a register and do any necessary sign extension or zero fill. It returns a string that can be used to reference the required register. As an example of the type-conversion process, given the following declarations:

Listing 6.90. op.c—Assignment Operators

```

571 value *assignment( op, dst, src )
572 int op;
573 value *dst, *src;
574 {
575     char *src_name;
576     char op_str[8], *p = op_str ;
577     value *val;
578
579     if( !dst->lvalue ) yyerror( "(=) lvalue required\n" );
580     if( !dst->is_tmp && dst->sym ) gen_comment( "%s", dst->sym->name );
581
582     /* Assemble the operator string for gen(). A leading @ is translated by
583      * gen() to a * at the far left of the output string. For example,
584      * ("@=",x,y) is output as "*x = y".
585      */
586
587     if( *dst->name != '&' ) *p++ = '@' ;
588     if( op ) *p++ = op ;
589     if( op == '<' || op == '>' ) *p++ = op ; /* <= or >= */
590     /* do always */ *p++ = '=' ;
591             *p++ = '\0' ;
592
593     src_name = rvalue( src );
594
595     if( IS_POINTER(dst->type) && IS_PTR_TYPE(src->type) )
596     {
597         if( op )
598             yyerror("Illegal operation (%c= on two pointers);\n", op);
599
600         else if( !the_same_type( dst->type->next, src->type->next, 0 ) )
601             yyerror("Illegal pointer assignment (type mismatch)\n");
602
603         else
604             gen( "=", dst->name + (*dst->name=='&' ? 1 : 0), src_name );
605     }
606     else
607     {
608         /* If the destination type is larger than the source type, perform an
609          * implicit cast (create a temporary of the correct type, otherwise
610          * just copy into the destination. convert_type() releases the source
611          * value.
612          */
613
614         if( !the_same_type( dst->type, src->type, 1 ) )
615         {
616             gen( op_str, dst->name + (*dst->name == '&' ? 1 : 0),
617                  convert_type( dst->type, src ) );
618         }
619         else
620         {
621             gen( op_str, dst->name + (*dst->name == '&' ? 1 : 0), src_name );
622             release_value( src );
623         }
624     }
625     return dst;
626 }

```

```
int i, j, k ;
long l ;
char *p;
```

simple input like `i=j=k;` generates the following output:

```
W(&_j) = W(&_k);
W(&_i) = W(&_j);
```

A more complex assignment like `*p=l=i` generates:

```
r0.w.low = W(&_i); /* convert i to long */
ext_word(r0);
L(&l) = r0.l; /* assign to l */
r0.l = L(&l); /* truncate l to char */
*BP(&p) = r0.b.b0; /* assign to *p */
```

The logical OR operator
(||).

The next level of *binary operators* handles the logical OR operator—the productions and related workhorse functions are in Listings 6.91, and 6.92.

Listing 6.91. c.y—Binary Operators: The Logical OR Operator and Auxillary Stack

```
143                                     /* Stacks. The stack macros are all in */
144                                     /* <tools/stack.h>, included earlier */
145 %
146 #include <tools/stack.h>          /* Stack macros. (see Appendix A) */
147
148 stk_err( o )
149 {
150     yyerror( o ? "Loop/switch nesting too deep or logical expr. too complex.\n"
151                 : "INTERNAL, label stack underflow.\n" );
152     exit( 1 );
153 }
154
155 #undef stack_err
156 #define stack_err(o)      stk_err(o)
157
158 stack_dcl (S_andor, int, 32); /* This stack wouldn't be necessary if I were */
159                                     /* willing to put a structure onto the value */
160                                     /* stack--or_list and and_list must both */
161                                     /* return 2 attributes; this stack will hold */
162                                     /* one of them. */
163 %}
164
165
166 or_expr : or_list           { int label;
167                               if( label = pop( S_andor ) )
168                                 $$ = gen_false_true( label, NULL );
169                               }
170
171     ;
172 or_list : or_list OROR    { if( $1 )
173                               or( $1, stack_item(S_andor,0) = tf_label() );
174                               }
175
176     and_expr        { or( $4, stack_item(S_andor,0) );
177                         $$ = NULL;
178                         }
179
180     | and_expr       { push( S_andor, 0 );
181                         }
182
183     ;
```

Listing 6.92. *op.c*—The Logical OR operator

```

627 void      or( val, label )
628     value   *val;
629     int      label;
630 {
631     val = gen_rvalue( val );
632
633     gen ( "NE",           val->name, "0"      );
634     gen ( "goto%s%d", L_TRUE,      label    );
635     release_value( val );
636 }
637 /*-----*/
638 value *gen_rvalue( val )
639 value *val;
640 {
641     /* This function is like rvalue(), except that emits code to generate a
642      * physical rvalue from a physical lvalue (instead of just messing with the
643      * name). It returns the 'value' structure for the new rvalue rather than a
644      * string.
645 */
646
647     if( !val->lvalue || *(val->name) == '&' ) /* rvalue or logical lvalue */
648         rvalue( val );                                /* just change the name */
649     else
650         val = tmp_gen( val->type, val );          /* actually do indirection */
651
652     return val;
653 }
```

The only difficulty here is the requirement that run-time evaluation of expressions containing logical OR operators must terminate as soon as truth can be determined. Looking at some generated output shows you what's going on. The expression *i*||*j*||*k* creates the following output:

```

NE(W(&_i),0)
  goto T1;
NE(W(&_j),0)
  goto T1;
NE(W(&_k),0)
  goto T1;
F1:
  W( T(1) ) = 0;
  goto E1;
T1:
  W( T(1) ) = 1;
E1:
```

The productions treat expressions involving `||` operators as an OR-operator-delimited list of subexpressions. The code that handles the individual list elements is on lines 724 to 730 of Listing 6.91. It uses the `or()` subroutine in Listing 6.92 to emit a test/branch instruction of the form:

```

NE(W(&_i),0)
  goto T1;
```

The code following the `F1` label is emitted on line 721 of Listing 6.91 after all the list elements have been processed.

The main implementation-related difficulty is that the *or_list* productions really want to return two attributes: the value that represents the operand and the numeric component of the label used as the target of the output **goto** statement. You could do this, of course, by adding a two-element structure to the value-stack union, but I was reluctant to make the value stack wider than necessary because the parser would slow down as a consequence—it would have to copy twice as much stuff on every parse cycle. I solved the problem by introducing an auxiliary stack (*S_andor*), declared in the **occ**s declaration section with the code on lines 146 to 163 of Listing 6.91. This extra stack holds the numeric component of the target label.

The best way to see what's going on is to follow the parse of *i || j || k* in Table 6.25. You should read the comments in that table now, looking at the source code where necessary to see what's going on.

Table 6.25. A Parse of *i || j || k*;

Parse Stack	<i>S_andor</i>	Comments
<i>stmt_list</i> □		Shift NAME . The shifted attribute is a pointer to the symbol that represents <i>i</i> .
<i>stmt_list</i> NAME □ <i>i</i>		Reduce by <i>unary</i> → NAME .
<i>stmt_list</i> unary □ <i>W(&_i)_L</i>		Reduce by <i>binary</i> → unary .
<i>stmt_list</i> binary □ <i>W(&_i)_L</i>		Reduce by <i>and_list</i> → binary .
<i>stmt_list</i> and_list □ <i>W(&_i)_L</i>		Reduce by <i>and_expr</i> → and_list .
<i>stmt_list</i> and_expr □ <i>W(&_i)_L</i>		Reduce by <i>or_list</i> → <i>and_list</i> . This is the first reduction in the list, so push 0 onto the <i>S_andor</i> stack. Note that, since this is the leftmost list element, the default <i>\$\$=\$\$1</i> is permitted to happen in order to tell the next reduction in the series what to do.
<i>stmt_list</i> or_list □ <i>W(&_i)_L</i>	0	Shift OROR .
<i>stmt_list</i> or_list OROR □ <i>W(&_i)_L</i> □	0	Reduce by the imbedded production in <i>or_list</i> → <i>or_list</i> {} <i>and_expr</i> . This is the first expression in the list—the compiler knows that it's first because <i>\$1</i> isn't NULL. Call <i>tf_label()</i> to get the numeric component of the target label, and then replace the 0 at the top of the <i>S_andor</i> stack with the the label number. This replacement tells the <i>or_expr</i> → <i>or_list</i> production that an has actually been processed—a reduction by <i>or_expr</i> → <i>or_list</i> occurs in every expression, even those that don't involve logical OR's. You can emit code only when the expression had an operator in it, however. The <i>S_andor</i> stack is empty if there wasn't one. The compiler emits the following code: CODE: <i>NE(W(&_i), 0)</i> CODE: <i>goto T1;</i>
<i>stmt_list</i> or_list OROR {128} □ <i>W(&_i)_L</i> □ □	1	Shift NAME . The shifted attribute is a pointer to the symbol that represents <i>j</i> .
<i>stmt_list</i> or_list OROR {128} NAME □ <i>W(&_i)_L</i> □ □ <i>j</i>	1	Reduce by <i>unary</i> → NAME . The synthesized attribute is an lvalue for <i>j</i> .
<i>stmt_list</i> or_list OROR {128} unary □ <i>W(&_i)_L</i> □ □ <i>W(&_j)_L</i>	1	Reduce by <i>binary</i> → unary .
<i>stmt_list</i> or_list OROR {128} binary □ <i>W(&_i)_L</i> □ □ <i>W(&_j)_L</i>	1	Reduce by <i>and_list</i> → binary .

continued...

The logical AND operator
(**&&**)

The code to handle the logical AND operator (**&&**) is almost identical to that for the OR operator. It's shown in Listings 6.93 and 6.94. The expression *i && j && k* generates the following output:

Table 6.25. Continued. A Parse of $i \mid\mid j \mid\mid k;$

Parse Stack					S_andor	Comments
stmt_list □ or_list W(&_i)_L	OROR {128}	and_list □ W(&_j)_L	1	Reduce by and_expr → and_list.		
stmt_list □ or_list W(&_i)_L	OROR {128}	and_expr □ W(&_j)_L	1	Reduce by or_list → or_list OROR {128} and_expr. Emit code to handle the second list element: CODE: NE(W(&_j), 0) CODE: goto T1; The numeric component of the label is at the top of the S_andor stack, and is examined with a stack_item() call. The post-reduction attribute attached to the or_list is NULL.		
stmt_list □ NULL	OROR	1	Shift OROR.			
stmt_list □ NULL	OROR	1	Reduce by imbedded production in or_list → or_list {} and_expr. This time, the attribute for \$1 is NULL, so no code is generated.			
stmt_list □ NULL	OROR {128}	1	Shift NAME. The shifted attribute is a pointer to the symbol that represents k.			
stmt_list □ NULL	OROR {128}	NAME □ k	1	Reduce by unary → NAME. The synthesized attribute is an lvalue for k.		
stmt_list □ NULL	OROR {128}	unary □ W(&_k)_L	1	Reduce by binary → unary.		
stmt_list □ NULL	OROR {128}	binary □ W(&_k)_L	1	Reduce by and_list → binary.		
stmt_list □ NULL	OROR {128}	and_list □ W(&_k)_L	1	Reduce by and_expr → and_list.		
stmt_list □ NULL	OROR {128}	and_expr □ W(&_k)_L	1	Reduce by or_list → or_list OROR {128} and_expr. Emit code to process the third list element. The numeric component of the label is at the top of the S_andor stack. CODE: NE(W(&_k), 0) CODE: goto T1; The synthesized attribute is also NULL, there.		
stmt_list □ NULL		1	Reduce by or_expr → or_list. The numeric component of the label is popped off the S_andor stack. If it's zero, then no operators were processed. It's 1, however, so emit the targets for all the goto branches generated by the previous list elements. CODE: F1: CODE: W(T(1)) = 0; CODE: goto E1; CODE: T1: CODE: W(T(1)) = 1; CODE: E1: The synthesized attribute is an rvalue for the temporary that holds the result of the OR operation.			
stmt_list □ W(T(1))_R						

```

EQ(W(&_i), 0) /* i && j && k */
  goto F1;
EQ(W(&_j), 0)
  goto F1;
EQ(W(&_k), 0)
  goto F1;
  goto T1;
F1:
  W( T(1) ) = 0;
  goto E1;
T1:
  W( T(1) ) = 1;
E1:

```

Since the run-time processing has to terminate as soon as a false expression is found, the test instruction is now an EQ rather than an NE, and the target is the false label rather than the true one. The same `S_andor` stack is used both for the AND and OR operators. The `goto T1` just above the `F1` label is needed to prevent the last list element from falling through to the false assignment. Note that the logical AND and OR operators nest correctly. The expression

```
(i || j && k || l)
```

generates the output in Listing 6.95 (`&&` is higher precedence than `||`). A blow-by-blow analysis of the parse of the previous expression is left as an exercise.

Listing 6.93. `c.y`—Binary Operators: The Logical AND Operator

```

733 and_expr: and_list      { int label;
734                               if( label = pop( S_andor ) )
735                               {
736                                   gen( "goto%s%d", L_TRUE, label );
737                                   $$ = gen_false_true( label, NULL );
738                               }
739                           }
740 ;
741 and_list: and_list ANDAND { if( $1 )
742                               and($1, stack_item(S_andor,0) = tf_label());
743                               }
744                               binary { and( $4, stack_item(S_andor,0) );
745                               $$ = NULL;
746                               }
747                               | binary { push( S_andor, 0 );
748                               }
749 ;
```

Listing 6.94. `op.c`—The Logical AND Operator

```

654 void and( val, label )
655 value *val;
656 int label;
657 {
658     val = gen_rvalue( val );
659
660     gen ( "EQ",      val->name, "0"      );
661     gen ( "goto%s%d", L_FALSE,   label   );
662     release_value( val );
663 }
```

Relational operators.

The remainder of the binary operators are handled by the *binary* productions, the first two right-hand sides of which are in Listing 6.96. The productions handle the relational operators, with the work done by the `relop()` subroutine in Listing 6.97. The `EQUOP` token matches either `==` or `!=`. A `RELOP` matches any of the following lexemes:

```
<=    >=    <    >
```

A single token can't be used for all six lexemes because the `EQUOPs` are higher precedence than the `RELOPs`. The associated, integer attributes are assigned as follows:

Listing 6.95. Output for `(i || j && k || l)`

```

1      NE(W(&_i),0)
2      goto T1;
3      EQ(W(&_j),0)
4      goto F2;
5      EQ(W(&_k),0)
6      goto F2;
7  F2:
8      W( T(1) ) = 0;
9      goto E2;
10 T2:
11     W( T(1) ) = 1;
12 E2:
13     NE(W( T(1) ),0)
14     goto T1;
15     NE(W(&_l),0)
16     goto T1;
17 F1:
18     W( T(1) ) = 0;
19     goto E1;
20 T1:
21     W( T(1) ) = 1;
22 E1:

```

Token	Lexeme	Attribute
EQUOP	==	'='
EQUOP	!=	'!'
RELOP	>	
RELOP	<	'<'
RELOP	>=	'G'
RELOP	<=	'L'

The `relop()` subroutine at the top of Listing 6.97 does all the work. It's passed the `relop()` operator's attribute and the two operands; `i < j` generates the following code:

```

LT( W(&_i), W(&_j) ) /* compare i and j. */
    goto T1;           /* Jump to true case on success; */
                      /* otherwise, fall through to false.*/
F1:   W( T(1) ) = 0;
      goto E1;
T1:   W( T(1) ) = 1;
E1:

```

The comparison directive on the first line (`LT`) changes with the incoming operand: `LT` for `<`, `GT` for `>`, `EQ` for `==`, and so on. `relop()` generates the code in several steps. Both operands are converted to `rvalues` by the `gen_rvalue()` calls on lines 673 and 674 of Listing 6.96. The `make_types_match` call on the next line applies the standard C type-conversion rules to the two operands, generating code to promote the smaller of the two variables to the type of the larger one. The subroutine starts on line 712 of Listing 6.96. It's passed pointers to the two values, and it might modify those values because the conversion might create another, larger temporary. Normally, `make_types_match` returns 1. It returns 0 if it can't do the conversion, as is the case when one operand is an `int` and the other a pointer.

`make_types_match()`

The **switch** on line 680 of Listing 6.96 translates the incoming attribute to an argument passed to the `gen()` call that generates the test on line 693. The **goto** is generated on the next line, and the numeric component of the target label is fetched from `tf_label()` at the same time. The code on lines 696 to 701 is doing a minor optimization. If the value passed into `gen_false_true()` is already a temporary of type **int**, then a second temporary is not created. If the original `v1` isn't an **int** temporary, the code on lines 696 to 701 swaps the two values in the hope that `v2` is a temporary.

Listing 6.96. c.y—Binary Operators: Relational Operators

```
750 binary      : binary RELOP      binary      { $$ = relop( $1, $2, $3 ); }
751          | binary EQUOP      binary      { $$ = relop( $1, $2, $3 ); }
```

Listing 6.97. op.c—Relational Operators

```
664 value    *relop( v1, op, v2 )
665 value    *v1;
666 int     op;
667 value    *v2;
668 {
669     char   *str_op ;
670     value  *tmp;
671     int    label;
672
673     v1 = gen_rvalue( v1 );
674     v2 = gen_rvalue( v2 );
675
676     if( !make_types_match( &v1, &v2 ) )
677         yyerror( "Illegal comparison of dissimilar types\n" );
678     else
679     {
680         switch( op )
681         {
682             case '>': /* > */    str_op = "GT";    break;
683             case '<': /* < */    str_op = "LT";    break;
684             case 'G': /* >= */   str_op = "GE";    break;
685             case 'L': /* <= */   str_op = "LE";    break;
686             case '!': /* != */   str_op = "NE";    break;
687             case '=': /* == */   str_op = "EQ";    break;
688             default:
689                 yyerror("INTERNAL, relop(): Bad request: %c\n", op );
690                 goto abort;
691         }
692
693         gen      ( str_op,      v1->name, v2->name           );
694         gen      ( "goto%s%d", L_TRUE,    label = tf_label()       );
695
696         if( !(v1->is_tmp && IS_INT( v1->type )) )
697         {
698             tmp = v1;                  /* try to make v1 an int temporary */
699             v1  = v2;
700             v2  = tmp;
701         }
702         v1 = gen_false_true( label, v1 );
703     }
704 }
```

Listing 6.97. continued...

```

705 abort:
706     release_value( v2 );                                /* discard the other value */
707     return v1;
708 }
709
710 /*-----*/
711
712 PRIVATE int make_types_match( v1p, v2p )
713 value **v1p, **v2p;
714 {
715     /* Takes care of type conversion. If the types are the same, do nothing;
716      * otherwise, apply the standard type-conversion rules to the smaller
717      * of the two operands. Return 1 on success or if the objects started out
718      * the same type. Return 0 (and don't do any conversions) if either operand
719      * is a pointer and the operands aren't the same type.
720      */
721
722     value *v1 = *v1p;
723     value *v2 = *v2p;
724
725     link *t1 = v1->type;
726     link *t2 = v2->type;
727
728     if( the_same_type(t1, t2, 0) && !IS_CHAR(t1) )
729         return 1;
730
731     if( IS_POINTER(t1) || IS_POINTER(t2) )
732         return 0;
733
734     if( IS_CHAR(t1) ) { v1 = tmp_gen(t1, v1); t1 = v1->type; }
735     if( IS_CHAR(t2) ) { v2 = tmp_gen(t2, v2); t2 = v2->type; }
736
737     if( IS ULONG(t1) && !IS ULONG(t2) )
738     {
739         if( IS LONG(t2) )
740             v2->type->UNSIGNED = 1;
741         else
742             v2 = tmp_gen( t1, v2 );
743     }
744     else if( !IS ULONG(t1) && IS ULONG(t2) )
745     {
746         if( IS LONG(t1) )
747             v1->type->UNSIGNED = 1;
748         else
749             v1 = tmp_gen( v2->type, v1 );
750     }
751     else if( IS LONG(t1) && !IS LONG(t2) ) v2 = tmp_gen( t1, v2 );
752     else if( !IS LONG(t1) && IS LONG(t2) ) v1 = tmp_gen( t2, v1 );
753     else if( IS_UINT(t1) && !IS_UINT(t2) ) v2->type->UNSIGNED = 1;
754     else if( !IS_UINT(t1) && IS_UINT(t2) ) v1->type->UNSIGNED = 1;
755
756     /* else they're both normal ints, do nothing */
757
758     *v1p = v1;
759     *v2p = v2;
760     return 1;
761 }
```

All other operators,
binary_op()

Most other binary operators are covered by the productions and code in Listings 6.98 and 6.99. Everything is covered but addition and subtraction, which require special handling because they can accept pointer operands. All the work is done in `binary_op()`, at the top of Listing 6.99. The routine is passed values for the two operands and an `int` that represents the operator. It generates the code that does the required operation, and returns a value that references the run-time result of the operation. This returned value is usually the incoming first argument, but it might not be if neither incoming value is a temporary. In addition, one or both of the incoming values is released.

Listing 6.98. c.y—Binary Operators: Other Arithmetic Operators

```

753 /* binary: */
754     | binary STAR      binary      { $$ = binary_op( $1, '*', $3 ); }
755     | binary DIVOP    binary      { $$ = binary_op( $1, $2, $3 ); }
756     | binary SHIFTOP   binary      { $$ = binary_op( $1, $2, $3 ); }
757     | binary AND       binary      { $$ = binary_op( $1, '&', $3 ); }
758     | binary XOR       binary      { $$ = binary_op( $1, '^', $3 ); }
759     | binary OR        binary      { $$ = binary_op( $1, '|', $3 ); }
```

Listing 6.99. op.c—Other Arithmetic Operators

```

762 value *binary_op( v1, op, v2 )
763 value *v1;
764 int op;
765 value *v2;
766 {
767     char *str_op;
768     int commutative = 0;           /* operator is commutative */
769
770     if( do_binary_const( &v1, op, &v2 ) )
771     {
772         release_value( v2 );
773         return v1;
774     }
775
776     v1 = gen_rvalue( v1 );
777     v2 = gen_rvalue( v2 );
778
779     if( !make_types_match( &v1, &v2 ) )
780         yyerror("%c%c: Illegal type conversion\n",
781                 (op=='>' || op=='<') ? op : ' ', op);
782     else
783     {
784         switch( op )
785         {
786             case '*':
787             case '&':
788             case '|':
789             case '^': commutative = 1;
790             case '/':
791             case '%':
792             case '<':                                /* << */ 
793             case '>':                                /* >> */ 
794                 dst_opt( &v1, &v2, commutative );
795         }
    }
```



Listing 6.99. continued...

```

796                     if( op == '<' )
797                         str_op = "<=";
798                     else if( op == '>' )
799                         str_op = IS_UNSIGNED(v1->type) ? ">L=" : ">>=";
800                     else
801                     {
802                         str_op = "X=";
803                         *str_op = op ;
804                     }
805
806                     gen( str_op, v1->name, v2->name );
807                     break;
808                 }
809             }
810             release_value( v2 );
811             return v1;
812         }
813
814 /*-----*/
815
816 PRIVATE int      do_binary_const( v1p, op, v2p )
817 value    **v1p;
818 int       op;
819 value    **v2p;
820 {
821     /* If both operands are constants, do the arithmetic. On exit, *v1p
822      * is modified to point at the longer of the two incoming types
823      * and the result will be in the last link of *v1p's type chain.
824      */
825
826     long   x;
827     link   *t1 = (*v1p)->type ;
828     link   *t2 = (*v2p)->type ;
829     value  *tmp;
830
831     /* Note that this code assumes that all fields in the union start at the
832      * same address.
833      */
834
835     if( IS_CONSTANT(t1) && IS_CONSTANT(t2) )
836     {
837         if( IS_INT(t1) && IS_INT(t2) )
838         {
839             switch( op )
840             {
841                 case '+': t1->V_INT += t2->V_INT;           break;
842                 case '-': t1->V_INT -= t2->V_INT;           break;
843                 case '*': t1->V_INT *= t2->V_INT;           break;
844                 case '&': t1->V_INT &= t2->V_INT;           break;
845                 case '|': t1->V_INT |= t2->V_INT;           break;
846                 case '^': t1->V_INT ^= t2->V_INT;           break;
847                 case '/': t1->V_INT /= t2->V_INT;           break;
848                 case '%': t1->V_INT %= t2->V_INT;           break;
849                 case '<': t1->V_INT <= t2->V_INT;           break;
850
851                 case '>': if( IS_UNSIGNED(t1) ) t1->V_UINT >>= t2->V_INT;
852                               t1->V_INT    >>= t2->V_INT;
853                               break;
854             }
855         }
856     }
857     return 1;

```



Listing 6.99. continued...

```

856         }
857     else if( IS_LONG(t1) && IS_LONG(t2) )
858     {
859         switch( op )
860         {
861             case '+':   t1->V_LONG += t2->V_LONG;           break;
862             case '-':   t1->V_LONG -= t2->V_LONG;           break;
863             case '*':   t1->V_LONG *= t2->V_LONG;           break;
864             case '&':   t1->V_LONG &= t2->V_LONG;           break;
865             case '|':   t1->V_LONG |= t2->V_LONG;           break;
866             case '^':   t1->V_LONG ^= t2->V_LONG;           break;
867             case '/':   t1->V_LONG /= t2->V_LONG;           break;
868             case '%':   t1->V_LONG %= t2->V_LONG;           break;
869             case '<':   t1->V_LONG <= t2->V_LONG;           break;
870
871         case '>':   if( IS_UNSIGNED(t1) ) t1->V ULONG >= t2->V_LONG;
872             else                      t1->V_LONG >= t2->V_LONG;
873             break;
874         }
875         return 1;
876     }
877     else if( IS_LONG(t1) && IS_INT(t2) )
878     {
879         switch( op )
880         {
881             case '+':   t1->V_LONG += t2->V_INT;           break;
882             case '-':   t1->V_LONG -= t2->V_INT;           break;
883             case '*':   t1->V_LONG *= t2->V_INT;           break;
884             case '&':   t1->V_LONG &= t2->V_INT;           break;
885             case '|':   t1->V_LONG |= t2->V_INT;           break;
886             case '^':   t1->V_LONG ^= t2->V_INT;           break;
887             case '/':   t1->V_LONG /= t2->V_INT;           break;
888             case '%':   t1->V_LONG %= t2->V_INT;           break;
889             case '<':   t1->V_LONG <= t2->V_INT;           break;
890
891         case '>':   if( IS_UNSIGNED(t1) ) t1->V ULONG >= t2->V_INT;
892             else                      t1->V_LONG >= t2->V_INT;
893             break;
894         }
895         return 1;
896     }
897     else if( IS_INT(t1) && IS_LONG(t2) )
898     {
899         /* Avoid commutativity problems by doing the arithmetic first,
900          * then swapping the operand values.
901         */
902
903         switch( op )
904         {
905             case '+':   x = t1->V_INT + t2->V_LONG;
906             case '-':   x = t1->V_INT - t2->V_LONG;
907             case '*':   x = t1->V_INT * t2->V_LONG;
908             case '&':   x = t1->V_INT & t2->V_LONG;
909             case '|':   x = t1->V_INT | t2->V_LONG;
910             case '^':   x = t1->V_INT ^ t2->V_LONG;
911             case '/':   x = t1->V_INT / t2->V_LONG;
912             case '%':   x = t1->V_INT % t2->V_LONG;
913             case '<':   x = t1->V_INT << t2->V_LONG;
914

```



Listing 6.99. continued...

```

915         case '>': if( IS_UINT(t1) ) x = t1->V_UINT >> t2->V_LONG;
916             else           x = t1->V_INT   >> t2->V_LONG;
917             break;
918         }
919
920         t2->V_LONG = x;      /* Modify v1 to point at the larger */
921         tmp = *v1p;          /* operand by swapping *v1p and *v2p. */
922         *v1p = *v2p;
923         *v2p = tmp;
924         return 1;
925     }
926 }
927 return 0;
928 }
929 /*-----*/
930
931 PRIVATE void dst_opt( leftp, rightp, commutative )
932 value **leftp;
933 value **rightp;
934 {
935     /* Optimizes various sources and destination as follows:
936     *
937     * operation is not commutative:
938     *     if *left is a temporary: do nothing
939     *     else:                   create a temporary and
940     *                           initialize it to *left,
941     *                           freeing *left
942     *                           *left = new temporary
943
944     * operation is commutative:
945     *     if *left is a temporary
946     *     else if *right is a temporary
947     *     else
948     */
949     value *tmp;
950
951     if( ! (*leftp)->is_tmp )
952     {
953         if( commutative && (*rightp)->is_tmp )
954         {
955             tmp = *leftp;
956             *leftp = *rightp;
957             *rightp = tmp;
958         }
959         else
960             *leftp = tmp_gen( (*leftp)->type, *leftp );
961     }
962 }

```

`binary_op()` starts out by trying to perform a type of optimization called *constant folding*. If both of the incoming values represent constants, then the arithmetic is done internally at compile time rather than generating code. The result is put into the last link of whichever of the two incoming values was larger, and that value is also the synthesized attribute. The work is done by `do_binary_constant()` starting on line 816 of Listing 6.99. An `if` clause is provided for each of the possible incoming types. Note that `do_binary_constant()` is passed pointers to the value pointers. The extra indirection is necessary because, if the left operand is larger than the right operand,

Constant folding.

the two values are swapped (after doing the arithmetic, of course). The code to do the swapping starts on line 920.

If constant folding couldn't be performed, then `binary_op` must generate some code. It starts out on lines 776 and 777 of Listing 6.99 by converting the incoming values to rvalues. It then does any necessary type conversions with the `make_types_match()` call on line 779. The `switch` on line 784 figures out if the operation is commutative, and the `dst_opt()` on line 794 juggles around the operands to make the code more efficient.

`dst_opt()`

`dst_opt()` starts on line 931 of Listing 6.99. It is also passed two pointers to value pointers, and it makes sure that the destination value is a temporary variable. If it's already a temporary, `dst_opt()` does nothing; otherwise, if the right operand is a temporary and the left one isn't, and if the operation is commutative, it swaps the two operands; otherwise, it generates a temporary to hold the result and copies the left operand into it.

Returning to `binary_op()`, the arithmetic instruction is finally generated on line 806 of Listing 6.99.

Addition and subtraction.

The last of the binary operators are the addition and subtraction operators, handled by the productions and code in Listings 6.100 and 6.101. The only real difference between the action here and the action for the operators we just looked at is that pointers are legal here. It's legal to subtract two pointers, subtract an integer from a pointer, or add an integer to a pointer. The extra code that handles pointers is on lines 1019 to 1057 of Listing 6.101.

The final group of expression productions are in Listing 6.102. They are pretty much self-explanatory.

Listing 6.100. *c.y*—Binary Operator Productions: Addition and Subtraction

```

760 /* binary: */
761     | binary PLUS      binary      { $$ = plus_minus( $1, '+', $3 ); }
762     | binary MINUS    binary      { $$ = plus_minus( $1, '-', $3 ); }
763     | unary
764 ;

```

Listing 6.101. *op.c*—Addition and Subtraction Processing

```

963 value *plus_minus( v1, op, v2 )
964 value *v1;
965 int op;
966 value *v2;
967 {
968     value *tmp;
969     int v1_is_ptr;
970     int v2_is_ptr;
971     char *scratch;
972     char *gen_op;
973
974     gen_op = (op == '+') ? "+=" : "-=";
975     v1 = gen_rvalue( v1 );
976     v2 = gen_rvalue( v2 );
977     v2_is_ptr = IS_POINTER(v2->type);
978     v1_is_ptr = IS_POINTER(v1->type);
979

```



Listing 6.101. continued...

```

980     /* First, get all the error checking out of the way and return if
981     * an error is detected.
982     */
983
984     if( v1_is_ptr && v2_is_ptr )
985     {
986         if( op == '+' || !the_same_type(v1->type, v2->type, 1) )
987         {
988             yyerror( "Illegal types (%c)\n", op );
989             release_value( v2 );
990             return v1;
991         }
992     }
993     else if( !v1_is_ptr && v2_is_ptr )
994     {
995         yyerror( "%c: left operand must be pointer", op );
996         release_value( v1 );
997         return v2;
998     }
999
1000    /* Now do the work. At this point one of the following cases exist:
1001    *
1002    *      v1:      op:      v2:
1003    *      number  [+/-]  number
1004    *      ptr     [+/-]  number
1005    *      ptr     -      ptr           (types must match)
1006    */
1007
1008    if( !(v1_is_ptr || v2_is_ptr) )          /* normal arithmetic */
1009    {
1010        if( !do_binary_const( &v1, op, &v2 ) )
1011        {
1012            make_types_match( &v1, &v2 );
1013            dst_opt( &v1, &v2, op == '+' );
1014            gen( gen_op, v1->name, v2->name );
1015        }
1016        release_value( v2 );
1017        return v1;
1018    }
1019    else
1020    {
1021        if( v1_is_ptr && v2_is_ptr )          /* ptr-ptr */
1022        {
1023            if( !v1->is_tmp )
1024                v1 = tmp_gen( v1->type, v1 );
1025
1026            gen( gen_op, v1->name, v2->name );
1027
1028            if( IS_AGGREGATE( v1->type->next ) )
1029                gen( "/=%s%d", v1->name, get_sizeof(v1->type->next) );
1030        }
1031        else if( !IS_AGGREGATE( v1->type->next ) )
1032        {
1033            /* ptr_to_nonaggregate [+/-] number */
1034            if( !v1->is_tmp )
1035                v1 = tmp_gen( v1->type, v1 );
1036
1037            gen( gen_op, v1->name, v2->name );
1038        }

```



Listing 6.101. continued...

```

1039         else                                /* ptr_to_aggregate [+-] number */
1040         {                                     /* do pointer arithmetic */
1041
1042             scratch = IS_LONG(v2->type) ? "r0.1" : "r0.w.low" ;
1043
1044             gen( "=",      "r1.pp", v1->name           );
1045             gen( "=",      scratch, v2->name           );
1046             gen( "*=%s%d", scratch, get_sizeof(v1->type->next) );
1047             gen( gen_op,   "r1.pp", scratch           );
1048
1049             if( !v1->is_tmp )
1050             {
1051                 tmp = tmp_create( v1->type, 0 );
1052                 release_value( v1 );
1053                 v1 = tmp;
1054             }
1055
1056             gen( "=", v1->name, "r1.pp" );
1057         }
1058     }
1059     release_value( v2 );
1060     return v1;
1061 }
```

Listing 6.102. c.y— High-Level Expression Processing

```

765 opt_expr
766     : expr          { release_value( $1 ); tmp_freeall(); }
767     | /* epsilon */
768     ;
769
770 const_expr
771     : expr          %prec COMMA
772     {
773         $$ = -1 ;
774
775         if( !IS_CONSTANT( $1->type ) )
776             yyerror("Constant required.");
777
778         else if( !IS_INT( $1->type ) )
779             yyerror("Constant expression must be int.");
780
781         else
782             $$ = $1->type->V_INT ;
783
784         release_value($1);
785         tmp_freeall();
786     }
787     ;
788
789 initializer : expr          %prec COMMA
790     | LC_initializer_list RC    { $$ = $2; }
791     ;
792
793 initializer_list
794     : initializer
795     | initializer_list COMMA initializer

```

Listing 6.102. continued...

```

796          {
797              yyerror("Aggregate initializers are not supported\n");
798              release_value( $3 );
799          }
800      ;

```

6.9 Statements and Control Flow

The only part of the compiler we've yet to examine is the statement productions, discussed in this section.

6.9.1 Simple Statements and **if/else**

It's best to start by looking at some examples. Table 6.26 shows input and output for a few simple control-flow statements. Figure 6.20 shows a more complex example of nested **if/else** statements. (I've shown the complete compiler output in the Figure.)

The productions and subroutines that generate this code are in Listings 6.103 and 6.104. The *stmt_list* production at the top just assembles a list of zero or more statements. There are no attributes. The simplest statement is defined on line 801 as a single semicolon. There is no action. A statement can also comprise a curly-brace-delimited compound statement (on line 811).

The next line defines a statement as an expression followed by a semicolon. The associated action frees the value holding the result of the expression evaluation and releases any temporary variables. Note that many expressions create unnecessary final values because there's no way for the parser to know whether or not an expression is part of a larger expression. For example, an assignment to a temporary is emitted as part of processing the the `++` operator in the statement:

```
a++;
```

but that temporary is never used. It is an easy matter for an optimizer to remove this extra assignment, which is, in any event, harmless.

The two forms of **return** statement are handled on lines 814 and 816 of Listing 6.103. The first is a simple return, with no value. The second takes care of the value by copying it into the required return-value register and then releasing the associated value and temporaries. Because returning from a subroutine involves stack-cleanup actions, a jump to a label immediately above the end-of-subroutine code is generated here rather than an explicit `ret()` instruction. The numeric part of the label is generated by `rlabel()`, in Listing 6.104. The end-of-subroutine code is generated during the reduction by

ext_def → *opt_specifiers funct_decl def_list compound_stmt*

(on line 563 of Listing 6.58, page 556), which executes an `rlabel(1)` to increment the numeric part of the label for the next subroutine.

Empty and compound statements.

Expression statements.

return statements.

Table 6.26. Simple Control-Flow: **return**, **goto** and **if/else**

Input	Output
return;	<pre> goto RET0; /* Generated by return statement */ ... RET0: /* Generated by end-of-subroutine */ unlink(); /* code. */ ret(); </pre>
return i+j;	<pre> W(T(1)) = W(&_i); /* compute i + j */ W(T(1)) += W(&_j); rF.w.low = W(T(1)); /* return value in register */ goto RET0: ... RET0: /* Generated in end-of-subroutine processing */ unlink(); ret(); </pre>
foo: ; goto foo;	<pre> _foo: goto _foo; </pre>
if(i < j) <i>++i;</i>	<pre> TST1: LT(W(&_i),W(&_j)) /* evaluate (i < j) and put */ goto T1; /* the result into T(1) */ F1: W(T(1)) = 0; goto E1; T1: W(T(1)) = 1; E1: EQ(W(T(1)),0) /* this test does loop control */ goto EXIT1; /* don't execute body if false */ W(&_i) += 1; /* body of the if statement */ W(T(1)) = W(&_i); EXIT1: </pre>
if(i < j) <i>++i;</i> else <i>++j;</i>	<pre> TST2: LT(W(&_i),W(&_j)) /* Evaluate (i < j) and put */ goto T2; /* the result into T(1). */ F2: W(T(1)) = 0; goto E2; T2: W(T(1)) = 1; E2: EQ(W(T(1)),0) /* This test does loop control. */ goto EXIT2; /* Jump to else clause if false.*/ W(&_i) += 1; /* Body of the if clause. */ W(T(1)) = W(&_i); goto EL2; /* Body of the else clause. */ W(&_j) += 1; W(T(1)) = W(&_j); EXIT2: EL2: </pre>

Figure 6.20. Nested **if/else** Statements

Input	Output
<pre> int i, j; fred() { int wilma; if(i) wilma=0; if(j) wilma=1; else if(j) wilma=2; else wilma=3; wilma=4; } </pre>	<pre> #include <tools/virtual.h> #define T(x) SEG(data) SEG(bss) common word _i; common word _j; #define L0 0 /* fred: locals */ #define L1 0 /* fred: temps. */ SEG(code) #undef T #define T(n) (fp-L0-(n*4)) PROC(_fred,public) link(L0+L1); /* fp-2 = x [variable] */ TST1: EQ(W(&_i),0) /* i */ goto EXIT1; W(fp-2) = 0; /* wilma */ TST2: EQ(W(&_j),0) /* j */ goto EXIT2; W(fp-2) = 1; /* wilma */ EXIT2: goto EL1; EXIT1: TST3: EQ(W(&_j),0) /* j */ goto EXIT3; W(fp-2) = 2; /* wilma */ EL3: W(fp-2) = 3; /* wilma */ W(fp-2) = 4; /* wilma */ EL1: unlink(); ret(); ENDP(_fred) </pre>

Listing 6.103. *c.y*—Statement Processing: **return**, **goto**, and **if/else**

```

801 stmt_list
802     : stmt_list statement
803     | /* epsilon */
804     ;
805
806 /*-----
807 * Statements
808 */
809
810 statement
811     : SEMI
812     | compound_stmt
813     | expr SEMI      { release_value($1); tmp_freeall(); }
814     | RETURN      SEMI { gen( "goto%s%d", L_RET, rlabel(0) ); }
815
816     | RETURN expr    SEMI { gen( "=", IS_INT ($2->type) ? "rF.w.low" :
817                               IS_POINTER($2->type) ? "rF.pp" :
818                               "rF.l",
819                               rvalue($2) );
820                               gen( "goto%s%d", L_RET, rlabel(0) );
821                               release_value( $2 );
822                               tmp_freeall();
823                         }
824     | GOTO target SEMI { gen("goto", $2); }
825     | target COLON   { gen(":", $1); }
826             statement
827
828     | IF LP test RP statement { gen( ":%s%d", L_NEXT, $3 ); }
829             }
830
831     | IF LP test RP statement ELSE { gen( "goto%s%d", L_ELSE, $3 );
832                                         gen( ":%s%d", L_NEXT, $3 );
833                                         }
834             statement { gen( ":%s%d", L_ELSE, $3 );
835                                         }
836

```

Listing 6.104. *op.c*—Get Numeric Part of End-of-Subroutine Label

```

1062 rlabel( incr )           /* Return the numeric component of the next */
1063 {                         /* return label, postincrementing it by one */
1064     static int num;         /* if incr is true. */
1065     return incr ? num++ : num;
1066 }

```

goto statements and labels.

The **goto** processing on lines 824 and 825 of Listing 6.103. is similarly straightforward. The compiler just generates jump instructions and labels as needed. The *target* nonterminal is at the top of Listing 6.105. It translates the label into a string which is returned as an attribute. Note that the action is imbedded in the middle of the label-processing production on line 825 of Listing 6.103. If the action were at the end, then the label would be generated after the associated statement was processed.

The next two productions, on lines 828 and 831 of Listing 6.103, handle the **if** and **if/else** statements. The best way to understand these actions is to compare the sample input and output in Table 6.26 and Figure 6.20. The code for the test is generated by the

test nonterminal, on line 843 of Listing 6.105, which outputs a label over the test code, generates the test code via *expr*, and then generates the statement that branches out of the test on failure. The first label isn't used here, but the same *test* production is used by the loop-processing productions, below, and these productions need a way to jump back up to the test code from the bottom of the loop. The presence of an extra label in an **if** statement is harmless.

Listing 6.105. *c.y*— Statement Processing: Tests and **goto** Targets

```

837 target  : NAME
838
839
840 }
841 ;
842
843 test   :
844
845 }
846     expr
847
848
849
850
851
852
853
854
855
856
857
858
859
860
861     /* empty */
862
863 ;
{   static char buf[ NAME_MAX ];
    sprintf(buf, "%.*s", NAME_MAX-2, yytext );
    $$ = buf;
}
{   static int label = 0;
gen( ":%s%d", L_TEST, $<num>$ = ++label );
}
{
    $$ = $<num>1;
    if( IS_INT_CONSTANT($2->type) )
    {
        if( ! $2->type->V_INT )
            yyerror("Test is always false\n");
    }
    else /* not an endless loop */
    {
        gen( "EQ",           rvalue($2), "0" );
        gen( "goto%s%d", L_NEXT,      $$ );
    }
    release_value( $2 );
    tmp_freeall();
}
{   $$ = 0; /* no test */
}
;
```

The returned attribute is the numeric component of all labels that are associated with the current statement. You'll need this information here to generate the target label for the exit branch, and the same numeric component is used to process the **else**. For example, the inner **if/else** at the bottom of Figure 6.20 uses three labels: TST3:, EXIT3, and EL3. The outer **if/else** uses TST1:, EXIT1, and EL1. The numeric component is generated in the *test* production, and the various alphabetic prefixes are defined symbolically in *label.h*. (It's in Listing 6.59 on page 557.)

Numeric component of label.

Finally, note that if the *expr* is a constant expression, no test is printed. An error message is printed if the expression evaluates to zero because the code in the body of the loop or **if** statement is unreachable; otherwise, the label is generated, but no test code is needed because the body is always executed. This way, input like:

Constant expressions in tests.

```
while( 1 )
    ...
```

evaluates to:

```
label:
...
goto label;
```

The alternative would be an explicit, run-time comparison of one to zero, but there's little point in that.

6.9.2 Loops, `break`, and `continue`

Loops.

Loops are handled by the productions in Listing 6.106; there's some sample input and output code in Table 6.27. The code generated for loops is very similar to that generated for an `if` statement. The main difference is a jump back up to the test at the bottom of the loop. The main difficulty with loops is `break` and `continue` statements, which are not syntactically attached to the loop-control productions. `break` and `continue` are treated just like labels by the productions on lines 915 and 923 of Listing 6.106. They can appear anywhere in a subroutine—there's nothing in the grammar that requires them to be inside a loop or `switch`. Nonetheless, you do need to know where to branch when a `break` or `continue` is encountered, and you need to detect a `break` or `continue` outside of a loop or `switch`.

The problem is solved with a few more auxiliary stacks, declared at the top of Listing 6.106. The top-of-stack item in `S_brk` is the numeric component of the target label for a `break` statement. The alphabetic component of the label is at the top of `S_brk_label`. I've used two stacks to save the trouble of calling `sprintf()` to assemble a physical label. The compiler pushes a label onto the stack as part of the initial loop-control processing (on line 866 of Listing 6.106, for example). It pops the label when the loop processing finishes on line 871 of Listing 6.106. `S_con` and `S_con_label` do the same thing for `continue` statements. If the stack is empty when a `break` or `continue` is encountered, the statement is outside of a loop, and a semantic error message is printed.

6.9.3 The `switch` Statement

The `switch` statement.

The final control-flow statement in the language is the `switch`. Switches can be processed in several different ways. First, bear in mind that a `switch` is really a vectored `goto` statement. Code like the following is legal, though it's bad style:

```
switch( i )
{
    case 0: if( condition )
              donald();
    else
    {
        case 1:   hewey();
                  dewie();
                  louie();
    }
    break;
}
```

You could do the same thing with `goto` statements as follows:

Table 6.27. Loops: **while**, **for**, and **do/while**

Input	Output
<pre>while(i < 10) { break; continue; }</pre>	<pre>TST3: LT(W(&<u>i</u>),10) /* Evaluate (i<10) and put */ /* the result into T(1). */ F3: W(T(1)) = 0; goto E3; T3: W(T(1)) = 1; E3: EQ(W(T(1)),0) /* Exit the loop if test fails.*/ /* Body of loop: */ /* goto EXIT3; */ goto EXIT3; /* break */ goto TST3; /* continue */ goto TST3; /* jump back up to the test */ EXIT3:</pre>
<pre>do { break; continue; } while(--i);</pre>	<pre>DTOP1: /* Top-of-loop marker */ /* Body of loop: */ goto DXIT1; /* break */ goto DTST1; /* continue */ DTST1: TST1: W(&<u>i</u>) == 1; /* Evaluate --i and put */ W(T(1)) = W(&<u>i</u>); /* the result into T(1). */ EQ(W(T(1)),0) /* Exit loop if test fails.*/ /* goto EXIT1; */ goto DTOP1; /* Jump back to top of loop */ DXIT1: EXIT1:</pre>
<pre>for(i = 0; i < 10; ++i) { break; continue; }</pre>	<pre>TST4: W(&<u>i</u>) = 0; /* Initialization part of for. */ /* Top-of-loop marker */ LT(W(&<u>i</u>),10) /* Evaluate (i<10) and put */ /* the result into T(1) */ F4: W(T(1)) = 0; goto E4; T4: W(T(1)) = 1; E4: EQ(W(T(1)),0) /* Exit the loop if test fails. */ /* goto EXIT4; */ goto BDY4; /* Skip over the increment */ /* /* increment portion of for stmt. */ INC4: W(&<u>i</u>) += 1; W(T(1)) = W(&<u>i</u>); goto TST4; /* Jump up to the test. */ /* /* Body of the loop */ goto EXIT4; /* /* break */ goto INC4; /* /* continue (jump to increment). */ goto INC4; /* Bottom of loop, jump to increment. */ EXIT4:</pre>

Listing 6.106. c.y—Statement Processing: Loops, **break**, and **continue**

```

164 %
165 /* These stacks are necessary because there's no syntactic connection break,
166 * continue, case, default and the affected loop-control statement.
167 */
168
169 stack_dcl (S_brk,      int,     32); /* number part of current break target */
170 stack_dcl (S_brk_label, char *, 32); /* string part of current break target */
171
172 stack_dcl (S_con,      int,     32); /* number part of current continue targ. */
173 stack_dcl (S_con_label, char *, 32); /* string part of current continue targ. */
174 %

864 /* statement: */
865
866     | WHILE LP test RP      { push(S_con, $3); push( S_con_label, L_TEST );
867                                         push(S_brk, $3); push( S_brk_label, L_NEXT );
868 }
869     statement            { gen( "goto%s%d", L_TEST, $3 );
870                               gen( ":%s%d",    L_NEXT, $3 );
871                               pop( S_con ); pop( S_con_label );
872                               pop( S_brk ); pop( S_brk_label );
873 }
874
875     | DO                  { static int label;
876
877             gen(":%s%d", L_DOTOP, $<num>$ = ++label );
878             push( S_con,           label );
879             push( S_con_label, L_DOTEST );
880             push( S_brk,           label );
881             push( S_brk_label, L_DOEXIT );
882 }
883     statement WHILE      { gen(":%s%d", L_DOTEST, $<num>2 );
884     LP test RP SEMI      { gen("goto%s%d", L_DOTOP, $<num>2 );
885                               gen( ":%s%d", L_DOEXIT, $<num>2 );
886                               gen( ":%s%d", L_NEXT,   $7 );
887                               pop( S_con );
888                               pop( S_con_label );
889                               pop( S_brk );
890                               pop( S_brk_label );
891 }
892
893     | FOR LP opt_expr SEMI
894         test      SEMI  {
895             gen("goto%s%d", L_BODY,      $5 );
896             gen(":%s%d",    L_INCREMENT, $5 );
897
898             push(S_con, $5);
899             push(S_con_label, L_INCREMENT );
900             push(S_brk, $5);
901             push(S_brk_label, L_NEXT );
902 }
903     opt_expr RP          { gen("goto%s%d", L_TEST,      $5 );
904                               gen(":%s%d",    L_BODY,      $5 );
905 }
906     statement            { gen("goto%s%d", L_INCREMENT, $5 );
907                               gen( ":%s%d",    L_NEXT,      $5 );
908 }
```



Listing 6.106. continued...

```

909                               pop( S_con );
910                               pop( S_con_label );
911                               pop( S_brk );
912                               pop( S_brk_label );
913 }
914 | BREAK SEMI      { if( stack_empty(S_brk) )
915                           yyerror("Nothing to break from\n");
916
917                           gen_comment("break");
918                           gen("goto%s%d", stack_item(S_brk_label,0),
919                                            stack_item(S_brk, 0));
920
921 }
922
923 | CONTINUE SEMI   { if( stack_empty(S_brk) )
924                           yyerror("Continue not in loop\n");
925
926                           gen_comment("continue");
927                           gen("goto%s%d", stack_item(S_con_label,0),
928                                            stack_item(S_con, 0));
929 }

```

```

if      ( i == 0 )  goto case0;
else if ( i == 2 )  goto casel;
else      goto end;

case0: if( condition )
        donald();
else
{
casel:   hewey();
        dewie();
        louie();
}
end:;

```

The simplest, and the least efficient, method of handling **switch** statements translates the **switch** and **cases** directly into a series of **if/else** statements. For example, code like this:

Translation to **if/else**.

```

switch( i )
{
case 0: /* code */
case 1: /* more code */
...
}

```

can be translated into the following C-code:

```

EQ( i, 0 )           /* test for case 0: */
    goto SW1;
    /* code */

    goto SW2;
SW1:                  /* Jump around test */
    EQ( i, 1 )
    goto SW3;
SW2:                  /* more code */
...

```

The main disadvantage of this method is run-time speed; a lot of unnecessary `goto` branches are needed to find the thing you're looking for. Moreover, you have to test every case explicitly to find the default one. In practice, translation to `if/else` is useful only if there are a very limited number of `case` statements.

One easy-to-do improvement eliminates the `goto` branches around the imbedded `case`-processing statements by moving all the tests to the bottom of the switch. This method is used by the current compiler, and is illustrated in Table 6.28. A jump at the top of the switch gets you to the test code, which then jumps back up to the correct place in the switch. A final `goto` branch jumps around the `case`-selection code if the last `case` doesn't have a `break` in it.

Table 6.28. Vectored Goto: `switch`

Input	Output
<pre> switch(i) { case 0: i = 0; /* fall through */ case 1: i = 1; break; default: i = 2; break; } </pre>	<pre> SW3: goto SW1; /* Jump to case-processing code.*/ /* Case 0 */ /* Fall through to next case. */ /* Case 1 */ /* break; */ /* default: */ /* break; */ /* Inserted by compiler in case */ /* last case has no break in it. */ /* Code to evaluate switch: */ EQ(W(&i),0) goto SW3; /* Jump to case 0. */ EQ(W(&i),1) goto SW4; /* Jump to case 1. */ goto SW5; /* Jump to default case. */ SW1: EQ(W(&i),0) goto SW3; /* Jump to case 0. */ EQ(W(&i),1) goto SW4; /* Jump to case 1. */ goto SW5; /* Jump to default case. */ SW2: </pre>

Dispatch tables.

The main reason that I'm using this approach is that limitations in C-code won't let me use the more efficient methods. It's worthwhile discussing these other methods, however. All of the alternate `switch` strategies use a table to compute the `goto` branches. The first method uses a data structure called a *dispatch table*—an array of two-member structures. The first member is the argument to the `case` statement, the second is a `goto` instruction that branches to the required location in the code. The switch:

```

switch( i );
{
case 0: washington();
case 1: jefferson();

```

```
case 3: adams();  
}
```

can be implemented with this table:

0	goto case0;
1	goto case1;
3	goto case3;

and this code:

```
case0: washington();  
case1: jefferson();  
case3: adams();
```

The compiler generates code to look up the case value in the table, often using a run-time subroutine call. If the value is found, the matching **goto** statement is executed, otherwise a branch to the default case is used. The table can be sorted by case value at compile time so the compiler can use a binary search at run time to find the proper label.

An even more efficient method uses a data structure called a *jump table*. A jump table is an array of **goto** vectors, indexed by case value. The compiler can process the switch with a single test and an array lookup:

```
if( argument to switch is in range )  
    goto jump_table[ argument_to_switch - value of smallest case ];  
else  
    goto default case;
```

The table elements are adjusted so that `jump_table[0]` corresponds to the smallest number used as an argument to a **case** statement. Holes in the table—elements for which there are no corresponding **case**—are filled with jumps to the default case. Note that, if the range of case values is less than or equal to twice the total number of case values, the jump table will be no larger than an equivalent dispatch table.

Though the current compiler can't use tables in the generated code, the method that it does use is easily adapted to this approach, because it assembles a compile-time model of the dispatch table. Listing 6.107 contains the data structures used for this purpose.

A dispatch-table element (`case_val`) is defined on lines three to seven. It contains two fields: `on_this` is the argument to the case statement, stored as an `int`, `go_here` is the numeric component of the label associated with this case; the alphabetic prefix is `SW`.

The `stab` structure, declared on line nine of Listing 6.107, is used to manage the table proper. It contains the following fields:

table	The dispatch table—an array of <code>case_val</code> structures.
cur	Points into <code>table</code> at the next available slot.
name	The name field of the <code>value</code> structure to which the expression argument of the switch statement evaluates. That is, a switch is recognized by the following production:

statement → **SWITCH LP** *expr RP compound_stmt*

The *expr* recognizes an entire expression—all code for evaluating the expression will have been generated when the *expr* is put onto the stack by the parser. This code creates a single `value` attribute that identifies the temporary variable that holds the run-time result of the expression evaluation; the `name` field from this `value` is copied to the `stab`'s `name` field.

Jump tables.

Jump table vs. dispatch table size.

def_label	The numeric component of the label associated with the default case. This number is 5 in the switch in Table 6.28.
stab_label	The numeric component of the labels that precede and follow the case-selection code. The label that follows the case-selection code has the value <code>stab_label+1</code> . This number is 1 in the switch in Table 6.28: SW1 precedes the selection code and SW2 follows it.

Listing 6.107. *switch.h*—Type Definitions for **switch** Processing

```

1 #define CASE_MAX 256           /* Maximum number of cases in a switch */
2
3 typedef struct case_val      /* a single dispatch-table element      */
4 {
5     int on_this;             /* The N in a "case N:" statement      */
6     int go_here;             /* Numeric component of label in output */
7 } case_val;                  /* code. */
8
9 typedef struct stab          /* a switch table */
10 {
11     case_val *cur;          /* pointer to next available slot in table */
12     case_val table[ CASE_MAX ]; /* switch table itself. */
13     char name [ VALNAME_MAX ]; /* switch on this rvalue */
14     int def_label;           /* label associated with default case */
15     int stab_label;          /* label at top and bottom of selector */
16                           /* code. Bottom label is stab_label+1. */
17 } stab;

```

The compiler keeps around a stack of pointers to stab structures, declared in Listing 6.108. The size of this stack limits the number of nested switch statements that are permitted in the input. A stab is allocated when the **switch** is processed, and a pointer to the stab is pushed. The pointer is popped after the entire switch has been processed. Every instance of a nested switch pushes a new stab pointer onto the stack, and that pointer is popped when the compiler is done with the nested switch.

Every **case** statement does two things: it emits a label of the form `SWn`, and it adds a new dispatch-table element for *n* to the stab structure at top of stack. (It modifies the entry pointed to by the `cur` field in the structure at top of stack, and then increments the `cur` field.) The numeric components of the label are just allocated sequentially by incrementing `Case_label`, also in Listing 6.108. The code that processes a **switch** is in Listings 6.108 and 6.109.

The only difficulty in the code is the point at which the value associated with the argument to the **switch** statement is released on line 945 of Listing 6.108. This is another compile-versus-run-time issue. The order of operations at run time is: (1) jump to the selector code at the bottom of the switch, (2) pick out the target label, (3) jump back up to that label. At run time, the value associated with the `expr` is used in step (2) to select the `goto` vector, and is not used further. Consequently, it can be recycled for use in the code that comprises the body of the switch. The selector code is not generated until after the body is processed, however. If the value was released after the selector code was printed, it could not be recycled in the **switch** body. Consequently, it's released at the top of the **switch**, before the body is processed. The value's name field is remembered in the stab structure so that the selector code can be generated later.

Listing 6.108. *c.y*— Statement Processing: The **switch** statement.

```
175  %{
176  int      Case_label = 0;           /* Label used to process case statements. */
177
178  stack_dcl (S_switch, stab *, 32); /* Switch table for current switch.      */
179  %}

930 /* statement: */
931     | SWITCH LP expr RP
932     {
933         /* Note that the end-of-switch label is the 2nd argument to
934         * new_stab + 1; This label should be used for breaks when in
935         * the switch.
936         */
937
938         push( S_switch, new_stab($3, ++Case_label) );
939         gen_comment("Jump to case-processing code" );
940         gen("goto%s%d", L_SWITCH, Case_label );
941
942         push( S_brk,          ++Case_label );
943         push( S_brk_label,    L_SWITCH      );
944
945         release_value( $3 );
946         tmp_freeall();
947     }
948     compound_stmt
949     {
950         gen_stab_and_free_table( pop( S_switch ) );
951     }
952
953     | CASE const_expr COLON
954     {
955         add_case   ( stack_item(S_switch,0), $2, ++Case_label );
956         gen_comment ( "case %d:", $2
957                     );
958         gen        ( ":%s%d" , L_SWITCH, Case_label );
959     }
960
961     | DEFAULT COLON
962     {
963         add_default_case( stack_item(S_switch,0), ++Case_label );
964         gen_comment("default:");
965         gen("::%s%d", L_SWITCH, Case_label );
966     }
;
```

Listing 6.109. *switch.c*—Switch Processing

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/debug.h>
4 #include <tools/hash.h>
5 #include <tools/l.h>
6 #include <tools/compiler.h>
7 #include <tools/c-code.h>
8
9 #include "symtab.h"
10 #include "value.h"
11 #include "proto.h"
12 #include "label.h"
13 #include "switch.h"
14
15 /*-----*/
16
17 PUBLIC stab *new_stab( val, label )
18 value *val;
19 int label;
20 {
21     /* Allocate a new switch table and return a pointer to it. Use free() to
22      * discard the table. Val is the value to switch on, it is converted to
23      * an rvalue, if necessary, and the name is stored.
24      */
25
26     stab *p;
27
28     if( !(p = (stab *) malloc(sizeof(stab)) ) )
29     {
30         yyerror("No memory for switch\n");
31         exit( 1 );
32     }
33
34     p->cur        = p->table ;
35     p->stab_label = label;
36     p->def_label   = 0;
37     strncpy( p->name, rvalue(val), sizeof(p->name) );
38     return p;
39 }
40
41 /*-----*/
42
43 PUBLIC void add_case( p, on_this, go_here )
44 stab *p;
45 int on_this;
46 int go_here;
47 {
48     /* Add a new case to the stab at top of stack. The 'cur' field identifies
49      * the next available slot in the dispatch table.
50      */
51
52     if( p->cur > &(p->table[CASE_MAX-1]) )
53         yyerror("Too many cases in switch\n");
54     else
55     {
56         p->cur->on_this = on_this ;
57         p->cur->go_here = go_here;
58         ++( p->cur );
59     }

```



Listing 6.109. continued...

```

60 }
61
62 /*-----*/
63
64 PUBLIC void add_default_case( p, go_here )
65 stab *p;
66 int go_here;
67 {
68     /* Add the default case to the current switch by remembering its label */
69
70     if( p->def_label )
71         yyerror("Only one default case permitted in switch\n");
72
73     p->def_label = go_here;
74 }
75
76 /*-----*/
77
78 PUBLIC void gen_stab_and_free_table( p )
79 stab *p;
80 {
81     /* Generate the selector code at the bottom of the switch. This routine is
82      * emitting what amounts to a series of if/else statements. It could just
83      * as easily emit a dispatch or jump table and the code necessary to process
84      * that table at run time, however.
85     */
86
87     case_val *cp;
88     char nbuf[20];
89
90     gen( "goto%s%d", L_SWITCH, p->stab_label + 1 );
91     gen( ":%s%d", L_SWITCH, p->stab_label );
92
93     for( cp = p->table ; cp < p->cur ; ++cp )
94     {
95         gen( "EQ%s%d", p->name, cp->on_this );
96         gen( "goto%s%d", L_SWITCH, cp->go_here );
97     }
98
99     if( p->def_label )
100         gen( "goto%s%d", L_SWITCH, p->def_label );
101
102     gen( ":%s%d", L_SWITCH, p->stab_label + 1 );
103     free( p );
104 }
```

6.10 Exercises

- 6.1. What are the consequences of changing the initial productions in the current C grammar to the following:

$$\begin{array}{lcl} \text{program} & \rightarrow & \text{ext_def_list} \\ \text{ext_def_list} & \rightarrow & \text{ext_def ext_def_list} \mid \epsilon \end{array}$$

Would the parser still work? For how long? This exercise points out the differences between theoretical and practical grammars.

- 6.2. Find all the bugs in the compiler presented in the current chapter and fix them.
- 6.3. Draw a system of symbol, link, and structdef structures that represents the following global-variable definition:

```
struct Montague
{
    long   * (*Romeo[2])[5];
    short  **Juliet[10];
    struct Montague *next;
}
* (* (* Capulet)[ 10 ])();
```

- 6.4. Modify the compiler in the current chapter as follows:
- Add 4-byte **floats** to the compiler.
 - Add 8-byte **doubles** to the compiler.
 - Add support for the **const**, **volatile**, and **register** storage classes.
 - Make structures generate lvalues: Add structure assignment and add the ability to pass a structure to a subroutine (and to return it) by value.
 - Add run-time initialization of automatic variables.
 - Add compile-time initialization of static-local and global aggregates.
 - Add support for signed and unsigned bit fields.
 - Add support for function prototypes. A hard error should occur when the types in a definition disagree with a prototype. A warning should be printed (and any required type conversions be performed) when the arguments of a call differ from those in the prototype.

You'll have to call run-time subroutines to do the floating-point operations because floating point is not supported in C-code.

- 6.5. The scope of a structure in the current chapter's compiler is from the point of declaration to the end of the input file. This is incorrect—structures should follow the same scoping rules as variables. Fix the problem.
- 6.6. Though the compiler in the current chapter behaves reasonably well when it's given legal input, it has difficulty when presented with syntactically correct, but semantically incorrect input. Semantic errors, such as `++` being applied to a non-lvalue or an attempt to multiply two pointers, should be handled by printing an error message, and then pretending that legal input had been encountered by putting a legitimate attribute onto the value stack instead of garbage. Improve the semantic-error recovery in the current compiler.
- 6.7. Most code generation has been concentrated into the `gen()` subroutine to facilitate portability to other target languages. Declarations are not handled by `gen`, however. Either modify `gen()` to handles declarations, or create a new `decl_gen()` subroutine that outputs all declarations. Modify the variable-declaration actions to use this subroutine.
- 6.8. The compiler currently generates the following code for `i && j && k`:

```

EQ(W(&_i),0)
  goto F1;
EQ(W(&_j),0)
  goto F1;
EQ(W(&_k),0)
  goto F1;
  goto T1;
F1:
  W( T(1) ) = 0;
  goto E1;
T1:
  W( T(1) ) = 1;
E1:

```

This code has an unreachable `goto` statement above the `F1` label. Modify the compiler to generate the following code for the earlier input:

```

EQ(W(&_i),0)
  goto F1;
EQ(W(&_j),0)
  goto F1;
NE(W(&_k),0)
  goto T1;
F1:
  W( T(1) ) = 0;
  goto E1;
T1:
  W( T(1) ) = 1;
E1:

```

- 6.9. (a) Redesign C-code so that it is a more appropriate language for communication with a back end: the sizes of all objects should all be 1, and the stack should be 1 element wide; Space for local variables should not be allocated in the compiler (information about them should be passed to the back end to do the allocation), and so forth. Modify the compiler to emit your modified code.
 (b) Write a back end that translates that code to a real assembly language.
- 6.10. The PostScript graphics-description language is a postfix assembly language used by many printers. It is described in Adobe System's *PostScript Language Reference Manual* (Reading, Mass.: Addison-Wesley, 1985). Write a C-to-PostScript Translator. You may define a subset of C for this purpose. All the standard PostScript functions should be supported as intrinsic functions (recognized as keywords by the compiler and processed directly).
- 6.11. The UNIX `troff` typesetting program is really an assembler. Its input is an assembly-language description of a document made up of interspersed text and “dot commands” (assembler directives). `Troff` translates that input into a typeset version of that document. Design a C-like language that can describe a document, and then implement a compiler that translates that language to `troff`. At minimum you must support `ints`, one-dimensional arrays, subroutines (which must be able to return `int` values), and `if/else`, `for`, and `while` statements. You should also support a basic `string` type, with operators for concatenation, comparison, and so forth. Finally, you should support an aliasing facility that lets you give reasonable names to built-in registers and special characters. You can use recursion to do the loops. For example, the following loop executes ten times:

```
.de Lp
.    if \\$1 \{\\
.        executing iteration \\$1
.        nr x \\$1-1
.        Lp \\nx \}
.
.\
.Lp 10      \" This call loops ten times
```

- 6.12. Write compilers that duplicate the UNIX **pic** and **eqn** preprocessors, translating descriptions of pictures and equations into **troff** commands.
- 6.13. Write an ANSI-C preprocessor. It must be able to handle the concatenation (##) and “stringizing” operator (#).
- 6.14. Write a program that reads a C input file and which outputs ANSI function prototypes for all functions in that file. Your program must be able to handle **#include** statements. Your output must be acceptable to a standard, ANSI compiler.
- 6.15. Write a program that reads a C file containing old-style function definitions like this:

```
apostles( mat, mark, luke, john, fred )
char *mat;
long mark;
double luke;
{
```

and which outputs C++-style definitions like this:

```
apostles( char *mat, long mark, double luke, int john, int fred )
{
```

Note that the default **int** type must be supplied for **john** and **fred**.

- 6.16. Modify the solution to the previous exercise so that it goes in the other direction, translating C++ definitions to old-style C definitions.
- 6.17. Modify the solution to the previous exercise to translate an ANSI-C input file to pre-ANSI, K&R C. If several input files are listed on the program’s command line, they should be translated individually. Output should be written to files with a *.k&r* extension. For example, *foo.c* should be translated, and the output placed in *foo.k&r*. (Use *.knr* on UNIX systems so that you don’t have to quote the name.)

At very least, function definitions, string concatenation, and function-name mapping [such as `remove()` to `unlink()`], should be handled. If the first eight characters of an identifier are not unique, the identifier should be replaced with a unique 8-character name. These replacements should carry from one input file to another in the case of nonstatic global variables and subroutine names. That is, if several C input files are specified on your translator’s command line, the files should be translated individually (not merged together), but the too-long global symbols that are translated in one file should be translated to the same arbitrary name in all files. Be careful to check that your arbitrary name is not already being used in the input file for some other purpose.

Structure assignment, passing structures to functions, and the returning of structures from a function should also be detected and translated into a form that can be handled by a compiler that supports none of these (you can use implicit `memcpy()` calls for this purpose).

- 6.18. Write a “beautifier” program that reads a randomly formatted C input file and which outputs the file with nice indenting showing the block level, lined up braces, and so forth. You should do intelligent things with comments, trying to line up the ‘/*’ and ‘*/’ tokens in end-of-line comments in neat columns, like this:

```
code          /* comment           */
more code    /* another comment      */
yet more code /* yet another comment */
```

If a comment won’t fit onto the end of a reformatted line, then it should be moved to the previous line. Multiple-line comments should be output as follows:

```
/* this is a
 * multiple-line
 * comment
 */
```

with the /* at the same indent level as the surrounding code.

- 6.19. One of the more useful features of C++ is “function overloading”. You are permitted to declare several instances of the same function, each of which takes arguments of different types, like this:

```
overload Sophocles;
int    Sophocles( long   Jocasta, short Oedipus ){...}
double Sophocles( double Jocasta, char *Oedipus ){...}
```

The overload keyword tells the compiler that multiple definitions of a function are permitted. The compiler determines which version of the function to call by examining the types of the actual arguments in a given call. If the types in a call don’t match the definition exactly, the standard type conversions are used to promote the argument to the correct type (a warning is printed if this is done). Do the following:

- (a) Modify the C compiler in the current chapter to support function overloading.
 - (b) Write a preprocessor that translates overloaded functions to a form that can be handled by a standard ANSI compiler.
- 6.20. The Awk Programming Language is discussed in [Aho2]. Write a compiler that converts an awk input file to C.
- 6.21. (a) Write a C-to-Pascal converter.
 (b) Write a Pascal-to-C converter. Your program must be able to handle nested subroutine declarations. The organization of the Pascal stack frame is discussed in Appendix B. This is a much harder problem than (a).
- 6.22. (a) Write a FORTRAN-to-C converter.
 (b) Write a C-to-FORTRAN converter. All C data types must be handled properly, especially translation of C pointers to FORTRAN array indexes. This is a much harder problem than (a).
- 6.23. Modify *virtual.h* so that the registers and run-time stack are both 16 bits wide rather than 32 bits wide. Eliminate the *lword* basic type and .*l* register selector, and redefine the *ptr* type to be 16 bits wide. Finally, modify the compiler so that this new stack and register width is supported. The C word widths must remain unchanged—the compiler should still support a 32-bit **long int**, 16-bit **int**, and 8-bit **char**. How could you modify the compiler to make this translation easier?
- 6.24. Add an autodecrement and autoincrement addressing mode to C-code indirect modes [*W(p++)*, *WP(--p)*, and so forth], then modify the grammar and associated code-generation actions so that the following inputs generate a minimum number of C-code instructions:

*p++ *++p *p-- *--p

- 6.25. Write a program that translates C-code into your favorite assembly language.
- 6.26. Write a C-code interpreter that takes a C-code program as input and executes that code directly. The following subroutines should be built into the interpreter (so that they can be called directly from a C-code program):

putb() Print the low byte of the top-of-stack item to standard output (in hex).
 putw() Print the low word of the top-of-stack item to standard output (in hex).
 putl() Print the long word at the top of stack to standard output (in hex).
 puti() Print the long word at the top of stack to standard output (in decimal).
 putc() Print the low byte of the top-of-stack item to standard output as an ASCII character.
 puts() Print the string in the array whose address is at the top of stack to standard output.
 getb() Input a hex byte from standard input and return it in rF.
 getw() Input a hex word from standard input and return it in rF.w.low.
 getl() Input a hex long word from standard input and return it in rF.
 geti() Input a decimal long word and put it in rF.
 getc() Input an ASCII character from standard input and return it in rF.b.b0.
 dump() Print the contents of all registers to standard output.
 dumps() Print the top 20 stack items to standard output.

Execution should begin with a subroutine called *main()*, which must be supplied in the source file. You may require this function to be the first (or last) one in the file if you wish.

- 6.27. If you haven't done so already, make the interpreter in the previous exercise into a window-based system. One window displays the code as it's being executed, a second displays the register and stack contents, a third shows the contents of selected static variables (entered by you at a command prompt), and a fourth shows the C input line that corresponds to the output being executed. Use line numbers that are inserted into the output by the lexical analyzer for this last function.
- 6.28. If you haven't done so already, modify the interpreter to support the following breakpoint types:
- Break when a specified source-code line is executed.
 - Break when a specified subroutine is called.
 - Break when a specified global variable is modified or takes on a specific value.
 - Break when a specified C-code instruction is executed.
- 6.29. If you haven't done so already, modify the interpreter to display the contents of both global and local variables by specifying their name at a command prompt. A local symbol should be displayed only when executing a subroutine that contains that symbol. You will need to pass symbol-table information between the compiler and interpreter to do this. Use an auxiliary file.

Optimization Strategies

This chapter looks at optimization strategies at a high level. It is not intended to be an in-depth discussion of the topic—something that would take a book of its own to cover adequately. The basic types of optimizations are described, however, and optimization techniques are discussed in a general way. The way that the choice of intermediate code affects optimization is also discussed.

Optimizations are easily divided into three categories: parser optimizations; linear, peephole optimizations; and structural optimizations. I'll look at these three categories one at a time.

7.1 Parser Optimizations

The first category includes all optimizations that can be done by the parser itself. The simplest of these really come under the category of generating good code to begin with: using logical lvalues rather than physical ones, minimizing the number of `goto` branches, and so forth. The other common parser optimization is *intrinsic-function generation*. Intrinsic function calls are translated directly by the compiler into code that does the action normally performed by the function. For example, `strcpy()` is often implemented as an intrinsic function. The lexeme `strcpy` is recognized by the compiler as a keyword, and there is a production of the form:

```
expr: STRCPY LP expr COMMA expr RP
```

in the grammar. The associated action generates code to copy one string to another rather than generating code to call the `strcpy()` function. Intrinsic functions are particularly useful with small workhorse functions, a call to which can often have a higher overhead than the code that does the work. Other common intrinsic functions are the math functions like `sin()`, `cos()`, and `sqrt()`.

Intrinsic functions.

Some of the optimizations discussed below can also be done directly by the parser (such as using a left shift to implement multiplication by a constant).

7.2 Linear (Peephole) Optimizations

The optimizations that can be done directly by the parser are, by necessity, limited in scope to single productions. It's difficult for the parser to optimize any code that takes more than one production to process. For example, the following compiler input:

```
int i;
...
i = 5;
++i;
return i + 1;
```

generates this output:

```
W(&_i)      = 5;
W(&_i)      += 1;
W( T(1) )   = W(&_i);
W( T(1) )   = W(&_i);
W( T(1) )   += 1;
rF.w.low    = W( T(1) );
ret();
```

The two identical assignments on the third and fourth lines are unavoidable because they are generated by two separate statements—the first by the `++i` and the second by the `i+1`. It's easy for a separate optimizer pass to recognize this situation, however, and eliminate the redundant assignment.

This section looks at various optimizations that cannot be done in the parser, but which can be done by an auxiliary optimizer pass that goes through the compiler's output in a linear fashion, from top to bottom. (It may have to go through the output several times, however.) This kind of optimizer is called a *peephole optimizer*—“peephole” because it examines small blocks of contiguous instructions, one block at a time, as if the code were being scrolled past a window and only the code visible in the window could be manipulated. The optimizer scans the code looking for patterns, and then makes simple replacements. Peephole optimizers are usually small, fast programs, needing little memory to operate.

The kind of optimizations that you intend to do is a major factor in deciding the type of intermediate language that the compiler should generate. A peephole optimizer is happiest working with triples or quads. I'll use the triples generated by the compiler in the last chapter for examples in the following discussion.

7.2.1 Strength Reduction

A *strength reduction* replaces an operation with a more efficient operation or series of operations that yield the same result in fewer machine clock cycles. For example, multiplication by a power of two can be replaced by a left shift, which executes faster on most machines. (`x*8` can be done with `x<<3`.) You can divide a positive number by a power of two with a right shift (`x/8` is `x>>3` if `x` is positive) and do a modulus division by a power of two with a bitwise AND (`x%8` is `x&7`).

Other strength reductions are less obvious. For example, multiplication by small numbers can be replaced by multiple additions: `t0*=3` can be replaced with

```
t1 = t0;
t0 += t1;
t0 += t1;
```

Combinations of shifts and additions can be used for multiplication by larger numbers: `t0*=9` can be replaced by:

```
t1 = t0;
t1 <= 3;
t0 += t1;
```

[That is, $t0 \times 9 \equiv (t0 \times 8) + t0 \equiv (t0 << 3) + t0$.] Larger numbers can also be handled this way: $t0^{*}=27$ can be replaced by:

```
t1 = t0;
t1 <= 1;
t0 += t1;
t1 <= 2;
t0 += t1;
t1 <= 1;
t0 += t1;
```

You can see what's going on by looking at how a binary arithmetic is performed. A binary multiplication is done just like a decimal multiplication. 27_{10} is 11011_2 , so the multiplication is done like this:

$$\begin{array}{r}
 \begin{array}{cccccccc} d & d & d & d & d & d & d & d \\ \times & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \end{array} \\
 \hline
 \begin{array}{c} d & d & d & d & d & d & d & d \\ d & d & d & d & d & d & d & d \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ d & d & d & d & d & d & d & d \\ d & d & d & d & d & d & d & d \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ + & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array} \\
 \hline
 \end{array}$$

If there's a *1* in the multiplier, then the multiplicand, shifted left by a number of bits corresponding to the position of the *1* in the multiplier, is added to the product. At worst, you need as many shift/add steps as there are bits in the denominator. It's up to the optimizer to determine at what point it takes longer to do the shifts and additions than it does to use a multiplication instruction, but on many machines the shift/multiply strategy is faster for all but very large numbers. This last optimization is a classic example of a code-size-versus-execution-speed trade-off. The faster addition-and-shift code is much larger than the equivalent multiplication instruction.

There are also nonarithmetic strength reductions. For example, many machines have several different forms of jump or **goto** instruction, and the optimizer can sometimes modify the code so that the more efficient instruction can be used.

Instruction optimization.

7.2.2 Constant Folding and Constant Propagation

We've already seen constant folding at work in the last chapter, because the parser itself can do it in a limited way. Put simply, the compiler itself does any arithmetic that involves constants, if it can. An expression like $a+2*3$ is treated like $(a+6)$.

Order of evaluation can prevent the parser from doing constant folding. For example, the parser can't optimize the following input because of the left-to-right order of evaluation:

```
a + 1 + 3
```

The parser processes the $a+1$ first and puts the result into a temporary variable. The three is then added to the temporary. The following code is generated:

```
W( T(1) )      = W(&_i);
W( T(1) )      += 1;
W( T(1) )      += 3;
```

It's easy for an independent optimizer to see that $T(1)$ is modified twice in succession

without being used between the modifications, however, and replace the foregoing with:

```
W( T(1) )      = W(&_i);
W( T(1) )      += 4;
```

Multiplication by one, addition and subtraction of zero, and shift by zero can be eliminated entirely.

Constant Propagation.

Constant folding is actually a simple case of a more general optimization called *constant propagation*. Many variables retain constant values over a large portion of their lifetimes. The compiler can note when a constant is assigned to a variable and use that constant rather than the variable when the variable is referenced in the code. The constant propagates until the variable is modified. For example, code like this:

```
_y = 5;
...
_x = _y;
```

can be replaced with

```
_y = 5;
...
_x = 5;
```

(Assignment of a constant to a variable is a more efficient operation than a memory-to-memory copy on most machines.) At a higher level, the loop in the following code:

```
int i, j;

for( i = 5, j = 0; j < i ; ++j )
    foo( i );
```

can be treated as if the compiler had seen the following:

```
for( i = 5, j = 0; j < 5 ; ++j )
    foo( 5 );
```

The optimizer is, in effect, keeping track of the contents of all variables that contain constants. It keeps a local symbol table with an entry for each active variable, and that internal entry is modified to track modifications in the code. The internal value is modified every time the variable is modified. For example, this code:

```
t0 = 1;
t0 += 5;
t1 = t0;
```

is translated to this:

```
t0 = 1;
t1 = 6;
```

The compiler initializes an internal copy of `t0` when the first line is encountered. It then modifies that internal copy by adding five to it when the `t0+=5` is encountered and discards the `t0+=5` instruction. Finally, the modified value is used when the substitution is made on the third input line.

Dead variables.

7.2.3 Dead Variables and Dead Code

One type of optimization often leads to another. In the previous example, `t0` can be discarded after the constant propagation because it is not used anywhere. `t0` is a *dead variable*—a variable that is initialized (and perhaps modified), but never referenced in a function call or to the right of an equals sign. The variable is considered dead from the last time it is used until it is reinitialized—like this:

```
t0 = _a;
t0 += 5;
_x = t0;
...
/* t0 is now dead */ 
t0 += 1; /* This instruction can be eliminated. */
...
t0 = _b; /* t0 is now resurrected. */
```

At a higher level, the `i` in the following code is a dead variable—all references to it can be eliminated:

```
foo(x)
{
    int i;
    i = x;
    ++i;
}
```

A *dead assignment* is an extreme case of a dead variable. A dead assignment is an assignment to a variable that is never used or modified. The initial assignment of `t0=_a` in the earlier example is not a dead assignment because `t0` is used later on in the code. Nonetheless, dead assignments are very common in the output from the compiler in the previous chapter because of the way that expressions are processed. All expressions evaluate to something, and there is often a copy into an rvalue as the last step of the expression evaluation. Code like this:

```
int x;
++x;
```

generates the following output:

```
W(fp-2)      += 1;          /* x */
W( T(1) )    = W(fp-2);
```

The final assignment to `T(1)` is a dead assignment because `T(1)` is never used. At a higher level, the earlier constant propagation example translated this code:

```
for( i = 5, j = 0; j < i ; ++j )
    ...
```

into this:

```
for( i = 5, j = 0; j < 5 ; ++j )
    ...
```

The `i=5` is now a dead assignment and can be eliminated.

Dead code is code that can't be reached or does nothing useful. For example, code generated by the following input can be removed entirely:

```
if( 0 )
    do_something();
```

This input is translated to the following C-code:

```
NE( 0, 0 )
goto label;
call( _do_something );
W(T(1)) = rF.w.low;
label:
```

It is optimized in two steps. First, the dead assignment on the fourth line can be eliminated, yielding:

Dead assignment.

```

        NE( 0, 0 )
        goto label;
    call( _do_something );
label:

```

Eliminating useless code.

Since the NE always evaluates true, all code between the following **goto** and the label can be eliminated because it's unreachable. Similar optimizations involve the elimination of useless code (like the

```

        goto F1;
F1:

```

which is generated at the end of a list of `&&` operators). The following, more complex example takes two passes to optimize:

```

t0 = 0;
NE(t0,0)
    goto label;
...
label:

```

The first pass folds the constant 0 into t0, yielding this:

```

NE(0,0)
    goto label;
...
label:

```

The test now always fails, so it and the following **goto** instruction can both be eliminated.

Had we started with the following code:

```

t0 = 1;
NE(t0,0)
    goto label;
...
label:

```

the constant would fold to:

```

NE(1,0)
    goto label;
...
label:

```

which always tests true. Consequently the NE, the **goto**, and all code up to the label are dead and can be discarded. Note that the dead-code elimination is affected if there is a second entry point into the block. Input like the following:

```

switch( hitter )
{
case eddie_murray:  x = 1;
    while( x )
    {
case tony_phillips:      ...
    }
}

```

(it's sick,¹ but it's legal) could generate output like this:

```
_x = 1;
_____
NE(W(&_x), 0)
    goto label      Remove only this code.
...
_____
label2:
...
label:
```

The constant 1 propagates into the test, resulting in dead code, but only the code up to label2 can be eliminated because the **switch** could branch to the interior label. A region of code that has a single entry point and a single exit point (everything from a label or PROC directive up to the next **goto** or **ret ()** instruction) is called a *basic block*, and some optimizations, such as the current one, can be performed only within the confines of a basic block.

As we just saw, dead variable and code elimination can interact with constant propagation in interesting ways. The *i* in the following code isn't a dead variable—at least initially:

```
foo()
{
    int i, j, array[10];
    i = 5;
    ++i;
    j = array[i];
    ...
}
```

The compiler can propagate the constant 5 through to the array access, however:

```
foo()
{
    int i, j, array[10];
    i = 5;
    ++i;
    j = array[6];
    ...
}
```

Having done that, the *i* is now dead and can be eliminated:

```
foo()
{
    int j, array[10];
    j = array[6];
    ...
}
```

A form of dead-assignment elimination also applies to nonconstants. Code like the following, which is very common compiler output:

Basic block.

Interaction between optimizations.

1. That's a technical term.

Nonconstant dead assignments.

```
t0 = _x;
...
_y = t0;
```

can be replaced by:

```
_y = _x;
```

unless _x is modified before t0 is used. For example, this code:

```
t0 = _x;
_y = t0;
t0 = _z;
```

can be optimized to:

```
_y = _x;
t0 = _z;
```

but the following cannot be optimized because _x is modified before t0 is used:

```
t0 = _x;
_x += 1;
_y = t0;
```

Hardware problems,
volatile.

All the foregoing optimizations can be a real headache when you're interfacing to hardware. For example, an 8-bit, memory-mapped I/O port at physical address 0x10 can be modeled as follows on many machines:²

```
char *port;
port = (char *)0x10;
```

Thereafter, you can access the port with `*port`. The following code, for example, is intended to pulse the low bit of the output port at periodic intervals:

```
*port = 0; /* initialize output port to all zeros */
while( *port ) /* read input port, terminate if data available */
{
    *port = 1; /* pulse the low bit of the output port */
    *port = 0;
    delay();
}
```

Unfortunately, many optimizers eliminate all of the previous code. The `*port=1` is a dead assignment because `*port` is modified again before the value 1 is used, so the first optimization results in the following code:

```
*port = 0;
while( *port )
{
    *port = 0;
    delay();
}
```

Next, since the value of `*port` never changes (the same constant is always assigned to it), constant propagation takes over and the following code is produced:

2. or, better yet, with:

```
char *const port = 0x10;
```

The **const** says that `port` itself won't change value. You can also use a macro:

```
#define port ((char *) (0x10))
```

```
*port = 0;
while( 0 )
    delay();
```

The initial assignment to `*port` can now be eliminated because `*port` isn't used anywhere, and the loop can be discarded because its body is unreachable.

The ANSI **volatile** keyword is used to suppress these optimizations. No assumptions are made as to the value of any variable that's declared volatile—the compiler assumes that the value can change at any time, even if no explicit code modifies it. You can suppress all optimizations in the previous example by declaring `port` as follows:

```
volatile char *port;
```

or better yet:

```
volatile char *const port = 0x10;
```

This declaration says that the object at which `port` points is liable to change without notice, but that `port` itself is a constant whose value never changes.³

Suppress optimization with the **volatile** keyword.

7.2.4 Peephole Optimization: An Example

I'll demonstrate the mechanics of peephole optimization with a simple example. The input:

```
i = 5;
++i;
return i + 1;
```

generates this output:

```
W(&_i)      = 5;
W(&_i)      += 1;
W( T(1) )   = W(&_i);
W( T(1) )   = W(&_i);
W( T(1) )   += 1;
rF.w.low    = W( T(1) );
ret();
```

This output has several inefficiencies built into it, all of which can be eliminated by a peephole optimizer. The optimizer must operate on a basic block: The current code has no labels into it, and a single exit point at the bottom, so it comprises a basic block.

The first optimizations to apply are constant folding and propagation. The approach is straightforward: Go through the compiler output performing those operations that you can. You'll need a symbol table, accessed by variable name, that holds the contents of that variable insofar as it can be determined.

The optimizer's first pass evaluates expressions, if possible, modifying the table's contents to reflect the computed value. It also substitutes any reference to a variable that holds a constant value with that constant. Here, the optimizer reads the first line, creates a symbol table entry for `_i` and initializes it to 5. It now reads the second line. There's already an entry for `i` in the table, and `_i` is being modified by a constant. Consequently, it increments the symbol-table entry from five to six, and modifies the code as follows:

3. The cast can be done like this:

```
#define port ((volatile char *)0x10)
```

```

W(&_i)      = 5;
W(&_i)      = 6;
W( T(1) )   = W(&_i);
W( T(1) )   = W(&_i);
W( T(1) )   += 1;
rF.w.low    = W( T(1) );

```

Now it reads the third and fourth lines. Seeing that `_i` holds a constant, it replaces the reference to `_i` with that constant as follows:

```

W(&_i)      = 5;
W(&_i)      = 6;
W( T(1) )   = 6;
W( T(1) )   = 6;
W( T(1) )   += 1;
rF.w.low    = W( T(1) );

```

It also makes a symbol-table entry for `T(1)`, initializing it to 6. Reading the fifth line, it sees a variable that holds a constant (`T(1)`) being modified by a constant, so adjusts the internal symbol table and the output as before, yielding:

```

W(&_i)      = 5;
W(&_i)      = 6;
W( T(1) )   = 6;
W( T(1) )   = 6;
W( T(1) )   = 7;
rF.w.low    = W( T(1) );

```

Reading the final line, since `T(1)` holds a constant, it can be replaced with its value:

```

W(&_i)      = 5;
W(&_i)      = 6;
W( T(1) )   = 6;
W( T(1) )   = 6;
W( T(1) )   = 7;
rF.w.low    = 7;

```

Though there are still the same number of operations as before, they've all been translated to simple assignments.

The output from the first pass is now simplified further by going through it a second time, eliminating all dead variables—variables that aren't used in the current block. The strategy is, again, straightforward: Clear the symbol table, then go through the code, creating entries for all variables that appear as a source operand. Make a second pass eliminating all variables whose destination operand is not in the table. Since `T(1)` is not used in the source fields of any of the triples, all assignments to it are eliminated, yielding:

```

W(&_i)      = 5;
W(&_i)      = 6;
rF.w.low    = 7;

```

The assignments to `_i` cannot be eliminated because the `_i` could be used later on in the code.

The final optimization eliminates dead assignments, starting with an empty symbol table and proceeding as follows:

- (1) If a variable in the symbol table is modified by some instruction other than assignment, remove it from the table.
- (2) When a variable appears as the destination in an assignment:
 - a. If it's not in the table, add it, remembering the line number of the assignment

instruction.

- b. If the variable is in the table, discard the line referenced in the symbol-table entry, and then replace that symbol-table reference with one for the current line.

In the current example, the first assignment to `_i` is a dead assignment because `_i` is reinitialized before being used, so it is eliminated:

```
W(&_i)      = 6;
rF.w.low = 7;
```

7.3 Structural Optimizations

So far, all the optimizations that we've looked at can be performed by analyzing the linear output stream of the compiler. Some optimizations must analyze the overall structure of the code, however, and more sophisticated techniques must be used. All of these optimizations need to know something about the structure of the code being optimized. Consequently, they must work on a parse or syntax tree that represents the code, not on a series of instructions. One solution to this problem is for the parser to build a physical parse tree with structures and pointers rather than generating code; the parse tree can then be optimized, and finally traversed depth first to do the code generation. There are two difficulties with this approach. First, the optimizer can get pretty big, and there is often insufficient room in memory for it to be combined with the parser—two independent programs are needed. Similarly, the more code that you can get into memory at once, the more of the program's structure can be seen. There's no point in wasting memory for a parser. Finally, the entire parse tree isn't needed for optimization—a syntax tree representing the various statements and expressions is sufficient. The usual solution to the problem is for the parser to generate an intermediate language from which the syntax tree can be reconstructed, and that intermediate code is processed by the optimizer.

7.3.1 Postfix and Syntax Trees

The intermediate language most appropriate for this purpose is *postfix* or *Reverse-Polish Notation* (RPN). User's of Hewlett-Packard calculators and UNIX's `dc` desk-calculator program will already be familiar with this representation. In postfix notation, operands are pushed onto a stack without modification. Operators access the top few items on the stack, replacing them with the result of the operation. For example, the following C fragment

```
A * B + C * D
```

can be evaluated using the following postfix operations:

```
push A
push B
pop two items, multiply them, push the result
push C
push D
pop two items, multiply them, push the result
pop two items, add them, push the result
```

Postfix, Reverse-Polish Notation (RPN).

The result is on the top of the stack when the evaluation finishes. The two multiplications have to be done before the addition because `*` has higher precedence than `+`.

A postfix intermediate language is easy to generate because the compiler doesn't have to worry about assigning temporary variables—it just uses the run-time stack as its

scratch space. Postfix is useful in interpreter implementations because postfix virtual machines are easy to implement. The interpreter can translate the source code into a postfix form, which is then executed in software on the virtual machine.

The main advantage of postfix, here, is that the optimizer can reconstruct the entire parse tree—or to be more exact, a compacted form of the parse tree called a *syntax tree*—from the list of instructions.

A common postfix representation uses one instruction per line. If that instruction is an rvalue reference, the contents of the variable are pushed onto the stack. If it's an lvalue reference, the address is pushed. If it's an arithmetic instruction, the top one or two stack elements are manipulated, and they are replaced by the result, which is either an lvalue or rvalue depending on the operator. The expression $A \cdot B + A \cdot B$ is represented as:

```
AL
BL
*
AL
BL
*
+
+
```

There's no need for an explicit “push” operator as long as the operators can be distinguished from variable names. Similarly, explicit parentheses are never necessary because the order of evaluation is determined by the sequence of operations, and all incoming variables are treated as lvalues because there's no need for explicit temporaries in the intermediate code. (Consequently, the subscripts in the earlier example are redundant and can be omitted.)

Generating postfix intermediate code.

The first order of business is making a parser that generates postfix intermediate code for expressions. I've done this in Listings 7.1 and 7.2. The code is simplified because no temporary variables are necessary.

Listing 7.1. *postfix.y*—Parser with Postfix Code Generation

```

1 %term ICON NAME
2 %left PLUS MINUS
3 %left TIMES DIVIDE
4 %%
5 s      : expr
6 ;
7
8 expr   : NAME          { yycode ("%s\n", yytext ); }
9 | ICON           { yycode ("%s\n", yytext ); }
10 | expr DIVIDE expr { yycode ("/\n" ); }
11 | expr TIMES expr  { yycode ("*\n" ); }
12 | expr PLUS expr   { yycode ("+\n" ); }
13 | expr MINUS expr  { yycode ("-\\n" ); }
14 ;
15 %%
16 main () { yyparse (); }
```

Converting the postfix intermediate code to a syntax tree.

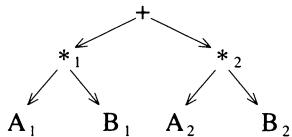
The next task is converting the intermediate code generated by our compiler back to a syntax tree. You can look at the generated intermediate code as a convenient way for the parser to pass the tree to the optimizer. The following syntax tree can be created when you run $A \cdot B + A \cdot B$ through the parser we just looked at:

Listing 7.2. *postfix.lex*— Lexical Analyzer for Postfix Code Generation

```

1  %
2  #include "yyout.h"
3  %
4  %%
5  [a-zA-Z_][a-zA-Z_0-9]*  return NAME;
6  [0-9]+                  return ICON;
7  "/"                     return DIVIDE;
8  "*"                     return TIMES;
9  "+"                     return PLUS;
10 "-"                    return MINUS;
11 .                      /* empty */
12 %%

```



All internal nodes represent operators and all the leaves reference lvalues. Also, the grouping of operators and operands is as you would expect, given the operator precedence built into the grammar.

Before demonstrating how to do this reconstruction, we'll need a data structure to represent the nodes in the tree. Listing 7.3 shows this data structure (a node) and a constructor subroutine that makes new nodes [`new()`]. The node structure is a normal binary-tree node, having left and right children. In addition, the name field holds variable names (A and B in this case) or the operator if the node is an internal node. The contents of this field will be modified by the optimizer, however. The op field usually holds the operator (`*` or `+`), but it is set to 0 in leaf nodes.

Data structures to represent the syntax tree: node.

Listing 7.3. *optimize.c*— The node Data Structure, Used to Construct Syntax Trees

```

1  #include <stdlib.h>           /* <malloc.h> for UNIX */
2
3  typedef struct node
4  {
5      char          name[16];
6      int           op;
7      struct node *left;
8      struct node *right;
9  }
10 node;
11
12 node *new()
13 {
14     node *p;
15     if( !(p = (node *) calloc( 1, sizeof(node) ) ) )
16         exit(1);
17     return p;
18 }

```

Reconstructing the syntax tree.

The `build()` subroutine in Listing 7.4 creates a syntax tree from a postfix input file. The input file must have one operand or operator per line and it must be perfect. That is, in order to simplify the code, I've dispensed with error detection. Input lines are read from standard input and the subroutine returns a pointer to the root node of the tree. The tree is built in a bottom-up fashion, using a local stack defined on line 22 to keep track of the partially-constructed tree. Figure 7.1 shows the syntax tree for the input discussed earlier as it is built—you'll notice the similarity between this process and the bottom-up parse process.

The `default` case on Line 30 is executed for variable names.⁴ It allocates and initializes a new node, and then pushes a pointer to the new node onto the stack. The child pointers are initialized to `NULL` by `new()`.

Operators are handled differently because they're internal nodes. A new node is allocated and initialized, then pointers to two existing nodes are popped and the child pointers of the new internal node are made to point at these. Finally, a pointer to the new node is pushed.

Listing 7.4. *optimize.c*—A Postfix to Syntax Tree Constructor

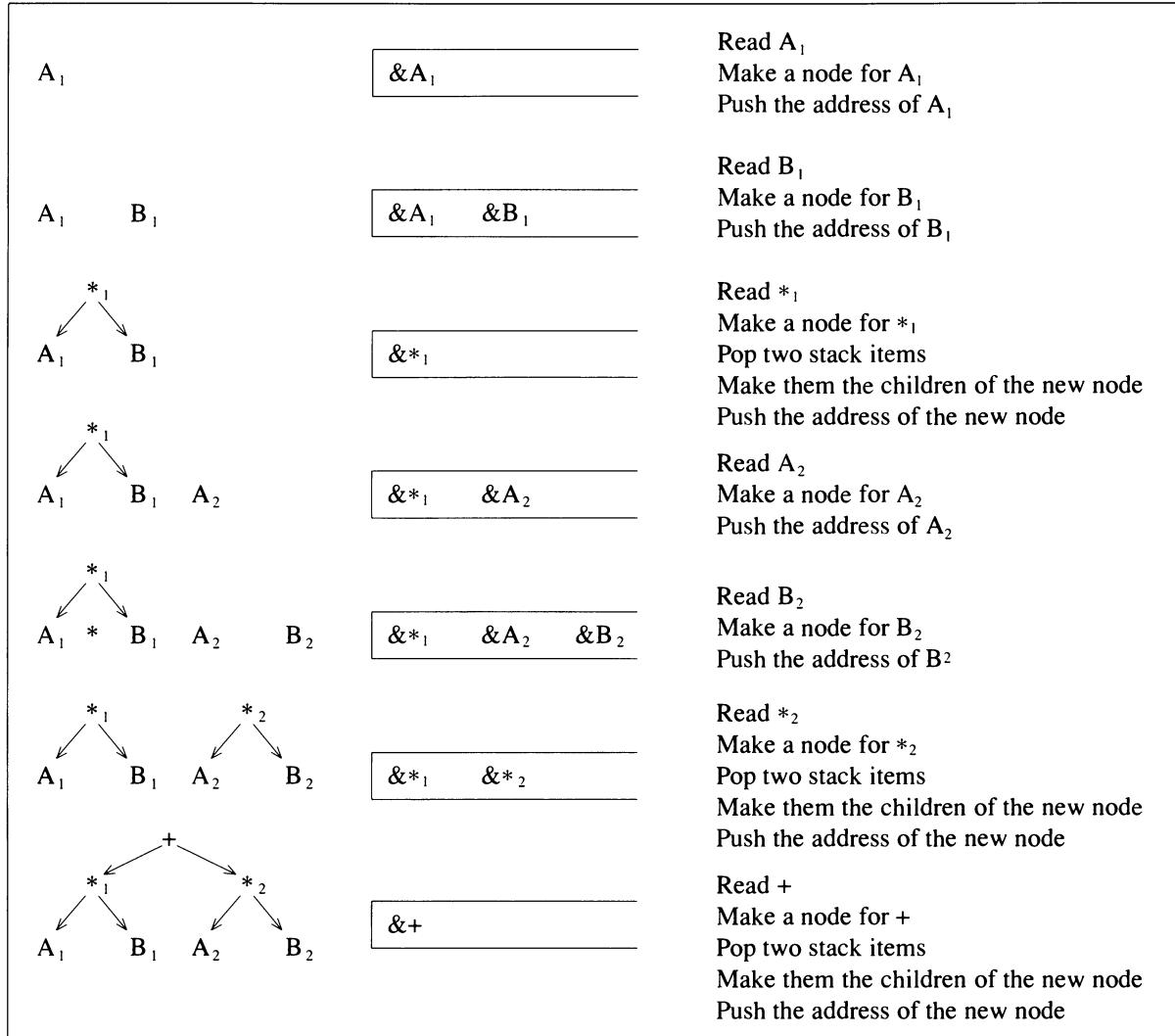
```

19  node    *build()
20 {
21     char    buf[80];
22     node   *stack[ 10 ];
23     node   **sp = stack - 1;
24     node   *p;
25
26     while( gets(buf) )
27     {
28         switch( *buf )
29         {
30             default: p = new();
31                     strcpy( p->name, buf );
32                     *++sp = p;
33                     break;
34
35             case '*':
36             case '+':
37                 p           = new();
38                 p->right   = *sp--;
39                 p->left    = *sp--;
40                 p->op      = *buf;
41                 p->name[0] = *buf;
42                 *++sp     = p;
43                 break;
44         }
45     }
46     return *sp--;
47 }
```

Generating code from the syntax tree.

Code can be generated from this tree by doing a depth-first traversal (visit the children then the parent). At every lvalue (ie. variable reference), generate an instruction of the form `temporary = variable`. At every internal node, generate the code

4. The `default` case can go anywhere in the switch. It doesn't have to be at the end.

Figure 7.1. Building a Syntax Tree

necessary to perform the operation on the temporaries that resulted from traversing the previous level, putting the result into a new temporary. The previously constructed tree generates the following output:

```
t0 = A
t1 = B
t1 *= t0
t2 = A
t3 = B
t3 *= t2
t3 += t1
```

The `trav()` subroutine in Listing 7.5 does the traversal. It takes the pointer returned from the previous `build()` call as its initial argument. If `root->op` is zero, then the current node is a leaf and you generate the code to move it to a temporary variable. The `sprintf()` call overwrites the `name` field with the name of the temporary variable. If the `op` field is nonnull, an interior node is being processed. In this case, `trav()` does an in-order traversal. The `if` statement is always true (for now—things

will change momentarily). The following `printf()` call prints the instruction, using the name fields of the two children to find out what temporaries to use. The `strcpy()` call then overwrites the name field of the current node to reflect the temporary that got the result of the last operation.

Listing 7.5. *optimize.c*— A Syntax-Tree Traversing, Code-Generation Pass

```

48 trav( root )
49 struct node *root;
50 {
51     static int tnum = 0;
52
53     if( !root )
54         return;
55
56     if( !root->op )      /* leaf */
57     {
58         printf ( "t%d = %s\n", tnum, root->name );
59         sprintf( root->name , "t%d", tnum );
60         ++tnum;
61     }
62     else
63     {
64         trav( root->left );
65
66         if( root->left != root->right )    /* Always true      */
67             trav( root->right );           /* unless optimized */
68
69         printf("%s %c= %s\n", root->right->name,
70                root->op, root->left->name );
71         strcpy( root->name, root->right->name );
72     }
73 }
```

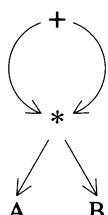
7.3.2 Common-Subexpression Elimination

The code that `trav()` outputs isn't too great, because the subexpression `A*B` is evaluated twice. It would be better to perform the multiplication only once and use the generated rvalue twice. You'd like the following output:

```

t0    =    A
t1    =    B
t1    *=   t0
t1    +=   t1
```

This transformation is called *common subexpression elimination*, and is a good example of the type of optimization that you can do by analyzing, and then modifying, the syntax tree. Since both subtrees of the `+` node are identical, the optimizer can eliminate one subtree and make both pointers in the `+` node point at the remaining subtree. The new syntax tree looks like this:



Both pointers in the + node point at the * node. This modified data structure is called a *Directed Acyclic Graph* or *DAG*. The DAG is created from the syntax tree by the `optimize()` function in Listing 7.6. This routine traverses the interior nodes of the tree, comparing the two subtrees. If the subtrees are identical, the left and right pointers of the parent node are made to point at the same child, effectively removing the other child from the tree. The comparison is done using the `makesig()` function, which traverses an entire subtree, assembling a string that shows the pre-order traversal (visit the root, the left subtree, then the right subtree) of the subtree by concatenating all the name fields. For example, the original syntax tree, when traversed from the root, creates the following signature string:

```
+**<A><B><A><B>
```

If two subtrees generate the same signature, they're equivalent.

Brackets are placed around the identifiers because the left and right subtrees in the following expression would incorrectly generate the same signatures if they weren't there:

```
(A * ABB) + (AA * BB)
```

Finally, you traverse the DAG using the `trav()` function that was developed earlier. That `if` statement now comes into play, preventing us from traversing the common subtree twice.

7.3.3 Register Allocation

Syntax trees and DAG's are useful for other optimizations as well. The simplest of these is efficient register allocation. In general, it's better to use registers than memory for temporary variables. It's difficult for the front end to do register allocation efficiently, however, because it doesn't know how many registers the target machine has. There are two problems with translating temporary-variable references to register references. First, you might need more temporary variables than there are registers, and the register allocation should be done in such a way that the registers are used as efficiently as possible: The temporaries that are used most frequently should be put into registers. The second issue is subroutine calls imbedded in expressions. If registers are used as temporaries, then they have to be pushed before calling the function and restored afterwards.

Both of these problems are easily solved by analyzing the parse or syntax tree for an expression before generating code for that expression. For example, an expression like this:

```
carl / philip + emanuel(a+b) * bach
```

generates the syntax tree in Figure 7.2 Each interior node in the tree represents the generation or modification of a temporary variable, and the temporary can appear at several interior nodes. The register-versus-nonregister problem is solved by assigning registers to those temporaries generated along the longest paths in the tree. Similarly, the optimizer can examine that tree and, noticing the function call, can decide to use registers only for those temporaries generated after the function returns or, as is the case with the function's argument, are pushed as part of the call.

Directed acyclic graph (DAG).

7.3.4 Lifetime Analysis

The next structure-related register optimization is *lifetime analysis*. This optimization takes care of a situation like the following:

Listing 7.6. *optimize.c*— A Subroutine for Common-Subexpression Elimination

```

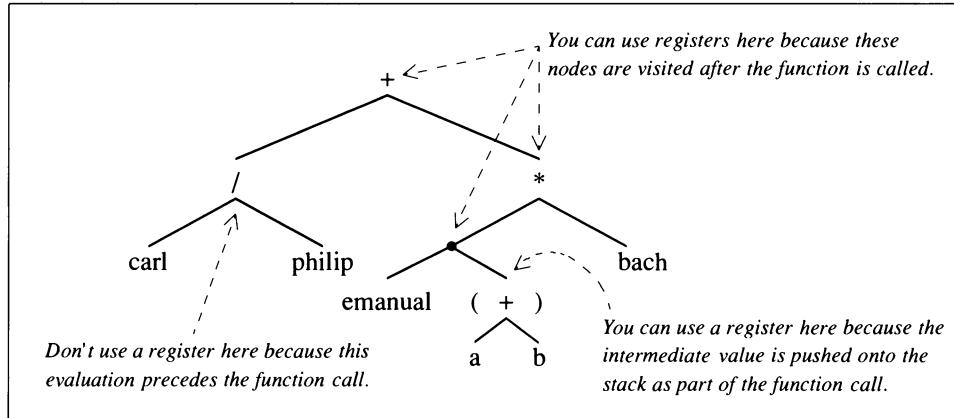
74 optimize( root ) /* Simplified optimizer--eliminates common subexpressions. */
75 node    *root;
76 {
77     char sig1[ 32 ];
78     char sig2[ 32 ];
79
80     if( root->right && root->left )
81     {
82         optimize( root->right );
83         optimize( root->left );
84
85         *sig1 = *sig2 = '\0';
86         makesig( root->right, sig1 );
87         makesig( root->left , sig2 );
88
89         if( strcmp( sig1, sig2 ) == 0 ) /* subtrees match */
90             root->right = root->left ;
91     }
92 }
93
94 makesig( root, str )
95 node    *root;
96 char   *str;
97 {
98     if( !root )
99         return;
100
101    if( isdigit( *root->name )
102        strcat( str, root->name );
103    else
104    {
105        strcat( "<", root->name );
106        strcat( str, root->name );
107        strcat( ">", root->name );
108    }
109    makesig( root->left, str );
110    makesig( root->right, str );
111 }
```

```

int i, j;

for( i = 10000; --i >= 0 ; )
    ...
for( j = 10000; --j >= 0 ; )
    ...
```

If the target machine has only a limited number of registers available for variables (say, one), then the **register** keyword can't be used effectively to improve the foregoing code. Ideally, because *i* is used only in the first loop and *j* only in the second, *i* should be placed in a register while the first loop is being executed, and *j* should be in a register while the second loop is executed. This determination can be made by analyzing the parse tree. In the simplest case, if a variable is found in only one subtree, then its lifetime is restricted to the code represented by that subtree. More sophisticated analysis can handle the more complex cases. For example, variables used for loop control can be given precedence over the variables in the loop body, unless the ones in the body are

Figure 7.2. Using the Syntax Tree for Register Allocation

used more frequently than the control variables.

7.3.5 Loop Unwinding

Loop unwinding is a method of optimizing loop execution speed, usually at the cost of code size. Put simply, the optimizer replaces the entire loop with the code that comprises the loop body, duplicated the number of times that the loop would execute at run time. This optimization, clearly, can be done only when the loop-control criteria can be computed at compile time. A loop like this:

```
for( i = 3; --i > 0 ; )
    foo(i);
```

can be replaced with the following code:

```
foo(3);
foo(2);
foo(1);
i = 0;
```

Something like this:

```
int array[10];
for( p = array; p <= array + 2; ++p )
    foo( p );
```

can be translated as follows:

```
r0.pp = array;
foo( r0.pp )
r0.pp += sizeof(int);
foo( r0.pp )
r0.pp += sizeof(int);
foo( r0.pp )
r0.pp += sizeof(int);
_p = WP( r0.pp );
```

7.3.6 Replacing Indexes with Pointers

If the values of an array index can be computed at compile time, not only can the loop be unwound, but the inefficient array indexes can be replaced by direct pointer operations. For example, this loop:

```
for( i = 0; i <= 2 ; ++i )
    array[i] = 0xff ;
```

can be translated as follows:

```
r0.pp      = array;
W(r0.pp)  = 0xff;
r0.pp      += sizeof(int);
W(r0.pp)  = 0xff;
r0.pp      += sizeof(int);
W(r0.pp)  = 0xff;
_i        = 3;
```

Note that the final assignment to *_i* is required in case *_i* is used further down in the code. Also, *_i* would have to be modified along with *r0.pp* if it were used in the loop body for something other than the array index.

7.3.7 Loop-Invariant Code Motion

Hoisting variables.

Loop-invariant optimizations analyze the body of a loop and *hoist* (move above the test portion of the loop) all of the code that doesn't change as the loop executes. For example, the division in the following loop is invariant:

```
for( i = 0; i < 10 ; ++i )
    array[i] += (numer/denom);
```

so you can modify the loop as follows:

```
r0.l = numer / denom ;

for( i = 0; i <= 2 ; ++i )
    array[i] += r0.l
```

The optimizations that we looked at earlier can also be applied here. Applying both array-index replacement and loop unwinding, you end up with the following C-code:

```
r0.l      = numer ;
r0.l      /= denom;
r1.pp    = array;
W(r1.pp) = r0.l ;
r1.pp    += sizeof(int);
W(r1.pp) = r0.l ;
r1.pp    += sizeof(int);
W(r1.pp) = r0.l ;
r1.pp    += sizeof(int);
i = 3;
```

Code-motion problems.

Note that hoisting division can cause problems if you're not careful. The following example bug is lifted from the Microsoft C (ver. 5.1) documentation. This loop:

```
for( i = 0; i <= 2 ; ++i )
    if( denom != 0 )
        array[i] += (numer/denom);
```

is optimized as follows:

```
r0.l = numer / denom ;

for( i = 0; i <= 2 ; ++i )
    if( denom != 0 )
        array[i] += r0.l ;
```

The division is moved outside the loop because it's invariant, but the preceding test is not moved with it. You'll now get a run-time divide-by-zero error if *denom* happens to

be zero.

You, as a compiler writer, must decide if it's worth the risk of doing this kind of optimization. It's difficult for the compiler to distinguish between the safe and dangerous cases, here. For example, many C compilers perform risky optimizations because the compiler writer has assumed that a C programmer can understand the problems and take steps to remedy them at the source code level. It's better to provide the maximum optimization, even if it's dangerous, than to be conservative at the cost of less efficient code. A Pascal programmer may not have the same level of sophistication as a C programmer, however, so the better choice in this situation might be to avoid the risky optimization entirely or to require a special command-line switch to enable the optimization.

7.3.8 Loop Induction

Loop induction is a sort of strength reduction applied to loop-control statements. The optimizer tries to eliminate the multiplication implicit in array indexing by replacing it with the addition of a constant to a pointer. For example, a loop induction on `a1` modifies this code

```
int a1[10][10], a2[10][10];

for( i = 0; i < 10; ++i )
    for( j = 0; j < 10; ++j )
        a1[i][j] = a2[i][j];
```

as follows:

```
t1 = array + 10;
for( t0 = a1; t0 < t1; t0 += 20 )      /* 20 == 10 * sizeof(int) */
{
    t3 = t0 + 20;                      /* 20 == 10 * sizeof(int) */

    for( t2 = t0; t2 < t3; t2 += sizeof(int) )
        *t2 = a2[i][j];

    j = 10;
}
i = 10;
```

7.4 Aliasing Problems

Many of the loop optimizations have problems with *pointer aliasing*—code in which a memory location can be accessed in more than one way (both by a pointer and directly, or by two pointers that contain the same address). For example, the following code creates an alias because both `x` and `*p` refer to the same memory location:

```
prometheus()
{
    int x, *p = &x;
}
```

Similarly, the following code creates less obvious aliases :

```

pyramus()
{
    int array[93];
    thisbe( array, array );
}

thisbe( p1, p2 )
int *p1, *p2 ;
{
    /* *p1 and *p2 alias one another, here */
}

int x;
pandora()
{
    int *p = &x;

    /* *p is an alias for x, here */
}

```

The problem here is that the loop optimizations might modify one of the aliased variables without modifying the other one—the compiler has no way to know that the two variables reference the same object.

7.5 Exercises

- 7.1. Write a subroutine that multiplies two arbitrary `ints` using nothing but shift and multiply operations. The result should be returned in a `long`.
- 7.2. (a) Modify the compiler in the last chapter to replace all multiplication by constants with shift/add sequences and to replace all division by constants with shift/subtract sequences.
 (b) Do the foregoing with a separate optimizer pass instead of modifying the compiler proper.
- 7.3. Write a peephole optimizer that works on the code generated by the compiler in the last chapter. At very least, it should do constant folding, dead-variable elimination, and dead assignment elimination.
- 7.4. Add an assignment operator (`=`) to the postfix tree constructor and optimizer presented earlier in this chapter. The input:

```

a
b
+
a
b
+
*
c
=

```

should generate the following:

```

t0 = a
t1 = b
t1 *= t0
t1 += t1
c = t1

```

Your program must be able to handle expressions like:

a = b = c

which should be translated to the following:

c
b
=

a
=

7.5. Modify your solution to the previous exercise to support the following C operators:

+	-	*	/	%	&		~	()
=	==	<	>	<=	>=	!	&&		
++	--								

Standard associativity and precedence should be used.

- 7.6. (a) Modify the compiler in the previous chapter to generate a postfix intermediate language instead of C-code.
(b) Write an optimizer pass that does common-subexpression elimination on that postfix intermediate code.
(c) Write a back end that traverses the syntax tree created by the common-subexpression optimizer and generates C-code.

Support Functions

This appendix contains a myriad small subroutines that are used throughout the rest of the book. Collectively, they comprise a library called *comp.lib*. A few of the routines in this library are described elsewhere (the input and table-printing functions in Chapter Two, for example), but I've concentrated the other routines in one place so that the rest of the book won't be cluttered with distracting references to small subroutines whose function is not directly applicable to the matter at hand. It's important that you are familiar with the routines in this appendix, however—at least with the calling conventions, if not with the code itself.

Since straight code descriptions can be pretty deadly reading, I've organized each section to start with a description of subroutine-calling conventions and a discussion of how to use the routines in that section. This way, you can skip the blow-by-blow code descriptions if you're not interested. Most of the code is commented enough so that additional descriptive text isn't necessary. There's no point in making an overly long appendix even longer with descriptions of things that are obvious by reading the code.

Of the subroutines in this appendix, you should be particularly familiar with the `<tools/debug.h>` file described in the next section and with the set and hash-table functions described in the sections after that. It's worthwhile to thoroughly understand this code. You need to be familiar only with the calling conventions for the remainder of the functions in this appendix, however. The sources are presented, but they're not mandatory reading.

Directories used.

The auxiliary files and subroutines described here are organized into several directories. Sources for subroutines of general utility are all in `/src/tools`, those that are more directly compiler related are in `/src/compiler/lib`. Most of the `#include` files used by the routines in this chapter are in `/include/tools`, and they are included with a `#include <tools/whatever.h>` (I've set up `/include` as my default library directory). The one exception is `curses.h`, which is in `/include` itself, so that you can use a UNIX-compatible `#include <curses.h>`.

A.1 Miscellaneous Include Files

Many *.h* files are included at the tops of the various *.c* files presented in this book. With the exception of the files discussed in this chapter, all files included with angle brackets (*<stdio.h>*, *<signal.h>*, and so forth) are ANSI-compatible, standard-library files and should be supplied by your compiler vendor (in this case, Microsoft). If they aren't, then your compiler doesn't conform to ANSI and you'll have to work a little harder to port the code presented here. Most of these standard *.h* files are also supported by UNIX—those that aren't are supplied on the distribution disk so that you can port the code to UNIX without difficulty.

A.1.1 *debug.h*—Miscellaneous Macros

The *debug.h* file is #included in virtually every *.c* file in this book. It contains several macro definitions that are of general utility, and sets up various definitions that are useful in debugging. It also contains several macros that take care of various 8086 portability issues. (It's useful, by the way, to look at the 8086-related macros even if you're not using an 8086—not only will you get a good idea of the things that you have to consider in order to make your code portable in general, but, since the 8086 is such a pervasive architecture, a specific knowledge of the machine lets you write your code so that it can be as portable as possible in any environment. It's a mistake to assume that your code will never have to move to the 8086 environment.)

Listing A.1. *debug.h*—Miscellaneous Macros of General Utility

```

1  #ifdef DEBUG
2  #define PRIVATE
3  #define D(x) x
4  #else
5  #define PRIVATE static
6  #define D(x)
7  #endif
8  #define PUBLIC
9
10 #ifdef MSDOS
11 #define MS(x) x
12 #define UX(x)
13 #define ANSI
14 #define _8086
15 #else
16 #define MS(x)
17 #define UX(x) x
18 #define O_BINARY 0      /* no binary input mode in UNIX open() */
19 #typedef long time_t; /* for the VAX, may have to change this */
20 #typedef unsigned size_t; /* for the VAX, may have to change this */
21 extern char *strupr(); /* You need to supply one. */
22#endif
23
24 #ifdef ANSI          /* If ANSI is defined, put arg lists into
25 #define P(x) x          /* function prototypes. */
26 #define VA_LIST ...     /* and use ellipsis if a variable number of args */
27 #else
28 #define P(x) ()         /* Otherwise, discard argument lists and translate */
29 #define void char        /* void keyword to int. */
30 #define VA_LIST _a_r_g_s /* don't use ellipsis */
31#endif
32

```



Listing A.1. continued...

```

33  /* SEG(p)          Evaluates to the segment portion of an 8086 address.
34  * OFF(p)          Evaluates to the offset portion of an 8086 address.
35  * PHYS(p)         Evaluates to a long holding a physical address
36  */
37
38 #ifdef _8086
39 #define SEG(p)   ( ((unsigned *)&(p))[1] )
40 #define OFF(p)    ( ((unsigned *)&(p))[0] )
41 #define PHYS(p)   (((unsigned long)OFF(p)) + ((unsigned long)SEG(p) << 4))
42 #else
43 #define PHYS(p)   (p)
44 #endif
45
46 /* NUMELE(array)    Evaluates to the array size in elements
47  * LASTELE(array)   Evaluates to a pointer to the last element
48  * INBOUNDS(array,p) Evaluates to true if p points into the array.
49  * RANGE(a,b,c)     Evaluates to true if a <= b <= c
50  * max(a,b)        Evaluates to a or b, whichever is larger
51  * min(a,b)        Evaluates to a or b, whichever is smaller
52  *
53  * NBITS(type)      Returns number of bits in a variable of the indicated
54  *                   type;
55  * MAXINT           Evaluates to the value of the largest signed integer
56  */
57
58 #define NUMELE(a)    (sizeof(a)/sizeof(*(a)))
59 #define LASTELE(a)   ((a) + (NUMELE(a)-1))
60 #define TOOHIGH(a,p) ((p) - (a) > (NUMELE(a) - 1))
61 #define TOOLOW(a,p)  ((p) - (a) < 0)
62 #define INBOUNDS(a,p) ( ! (TOOHIGH(a,p) || TOOLOW(a,p)) )
63
64 #define _IS(t,x)   (((t)1 << (x))!=0) /* Evaluate true if the width of a      */
65                                /* variable of type of t is < x. The !=0 */ 
66                                /* assures that the answer is 1 or 0      */
67
68 #define NBITS(t)   (4 * (1 + _IS(t, 4) + _IS(t, 8) + _IS(t,12) + _IS(t,16) \
69                      + _IS(t,20) + _IS(t,24) + _IS(t,28) + _IS(t,32) ) )
70
71 #define MAXINT (((unsigned)^0) >> 1)
72
73 #ifndef max
74 #  define max(a,b) ( ((a) > (b)) ? (a) : (b) )
75 #endif
76 #ifndef min
77 #  define min(a,b) ( ((a) < (b)) ? (a) : (b) )
78 #endif
79 #define RANGE(a,b,c) ( (a) <= (b) && (b) <= (c) )

```

DEBUG

Local objects: PRIVATE.

The debugging versions of several macros are activated when DEBUG is **#defined** (either with an explicit **#define** or, more commonly, with a **-DDEBUG** on the compiler's command line) before `<tools/debug.h>` is included. The PRIVATE definition on line two expands to an empty string in this case—it effectively disappears from the input. When DEBUG is not defined, the alternative definition on line five is activated, so all PRIVATE variables become **static** variables when you're not debugging. This mechanism lets you put normally invisible variables and subroutines into the link map—so that you know where they are when you're debugging—and then limit their scope to the current file when debugging is completed. The complement to PRIVATE is

PUBLIC, defined on line eight. It is provided for documentation purposes only—PUBLIC variables and subroutines are those whose scope is not limited.

The D(x) macro on lines three and six is used for printing debugging diagnostics. When DEBUG is defined, the macro expands to its argument, so

```
D( printf("a diagnostic\n"); )
```

expands to

```
printf("a diagnostic\n");
```

Note that the semicolon is inside the parentheses. When DEBUG is not defined, the macro expands to an empty string and the argument is discarded. The earlier "printf()" statement goes away. This method is preferable to the more usual

```
#ifdef DEBUG
printf("a diagnostic\n");
#endif
```

for several reasons: it's more readable; it requires only one line, rather than three; it doesn't mess up your indentation (many compilers require the # to be in the leftmost column). A multiple-line argument to D(x) must be done with backslashes, like this:

```
D(printf("%s in de %s wit' dinah, strumn' on de ol' %s", \
         someone, "kitchen", "banjo" ); )
```

There's a potential problem here in that commas that are part of the argument to D() can be confused with the comma that separates macro arguments. ANSI says that commas enclosed by parentheses are not to be treated as argument separators (which is why the previous definition worked). Your compiler may not handle this case correctly, however. In any event, you'll have to parenthesize expressions that use the comma operator to pass them to D():

```
D( (statement1,statement2); )
```

It's better style to use two D() invocations, however:

```
D( statement1; )
D( statement2; )
```

The macros on lines ten to 22 help you write portable code. The arguments to MS() are incorporated into the program only if MSDOS is **#defined** (it's defined automatically by the Microsoft compiler). Arguments to UX() (for UNIX) are active only if MSDOS is missing. I've also defined ANSI and _8086 macros that are active only if MSDOS is also active. Move the ANSI definition outside of the **#ifdef** if you've an ANSI-compatible UNIX compiler.

The P macro on lines 25 and 28 handles another ANSI-related portability problem.¹ Many compilers (including many UNIX compilers) can't handle function prototypes. This macro uses a mechanism similar to the D() macro discussed earlier to translate prototypes into simple **extern** declarations if ANSI is not defined. For example, given the following input:

```
int dimitri P(( int x, long y ));
```

if ANSI is defined, the P(x) macro evaluates to its argument and the following

Global objects: PUBLIC.

Debugging diagnostics,
D().

Commas in the D()
macro argument.

UNIX- and MS-DOS-specific
code: MS(), UX().

ANSI, _8086.

Transforming function
prototypes to simple **extern**, P().

1. This macro is taken from [Jaeschke], p. 142, and is used in the Whitesmiths compiler.

translation results:

```
int dimitri ( int x, long y );
```

otherwise, the macro discards its argument and evaluates to (), so the following will be created:

```
int dimitri ();
```

Getting rid of the ellipsis in the UNIX environment, VA_LIST.

The ARGS macro on lines 26 and 30 takes care of yet another problem with non-ANSI compilers. An ANSI subroutine with a variable number of arguments is defined as follows:

```
printf( format, ... )
char *format;
```

The trailing ellipsis means “followed by any number of arguments of an arbitrary type.” Earlier compilers don’t support this mechanism. Consequently, you’d like to replace the ellipsis with a dummy argument name. You can do this with *debug.h* using the ARGS macro as follows:

```
printf( format, VA_LIST )
char *format;
```

VA_LIST expands to an ellipsis when ANSI is defined, otherwise it expands to the dummy argument name _a_r_g_s.

8086-related macros, SEG(), OFF(), PHYS().

The macros on lines 39 to 41 of Listing A.1 handle portability problems inherent in the 8086 architecture. The SEG() and OFF() definition are ignored in other systems, and the PHYS() macro just evaluates to its argument in non-8086 environments. Before looking at the definitions for these macros, a short discussion of the 8086 internal architecture is necessary. The 8086 is really a 16-bit machine in that all internal registers are 16 bits wide. The machine has a 20-bit address space, however. The discrepancy is made up by using two registers to create a physical address. An address is formed from two 16-bit components representing the contents of the two registers: the *segment* and the *offset*. These are usually represented as two 4-digit hex numbers separated by a colon: 1234:5678. The segment portion is on the left, the offset on the right. A physical address is formed in the machine using the following formula:

```
physical_address == ( segment << 4 ) + offset;
```

Three special registers called CS, DS, SS hold the segment components of the addresses of objects in the code, data, and stack regions of your program. (A fourth segment register, called ES, is used if you need to access an object not in one of these segments).

Problems with segment:offset addressing.

This architecture has several ramifications. First, there are two pointer sizes. If the segment component doesn’t change over the life of the program, only the offset needs to be stored; pointers are 16 bits wide in this situation. For example, the DS register doesn’t have to change if there’s less than 64K of data and if the DS register is initialized to the base address of that 64K region. All cells within that region can be accessed using the fixed DS register by loading a 16-bit offset into the appropriate offset register—the segment portion of the address can be loaded once when the program loads. If the data area is larger than 64K, then 32-bit pointers composed of a 16-bit segment and a 16-bit offset component must be used. You must load both the segment and offset components into their respective registers to access the target variable. Typically, the segment is stored in the top 16 bits of the pointer and the offset is stored in the bottom 16 bits. The physical address is not stored. A 32-bit pointer that has both a segment and offset component is called a *far* pointer; a 16-bit pointer containing the offset component only is called a *near* pointer. Note that a program can have 16-bit data pointers and 32-bit function pointers and vice versa, depending on how much code or data there is.

Near and far pointers.

One of the other problems with the segmented architecture is that there are 4,096 ways to address every physical cell in memory. For example, the following all address physical address 0x10000: 1000:0000, 0FFF:0010, OFFE:0020, OFFD:0030, and so forth.

The PHYS(p) macro on line 41 of Listing A.1 evaluates to the physical address associated with an 8086 32-bit, far pointer—a pointer in which both the segment and offset components of the address are stored. It computes this address using the same procedure used by the 8086 itself, shifting the segment component left by four bits and adding it to the offset. The segment and offset portions are extracted by the SEG and OFF macros on the previous two lines, reproduced here:

Finding a physical address on an 8086,
PHYS().

Isolating segment and offset components,
SEG(), OFF().

```
#define SEG(p)    ( ((unsigned *) &(p)) [1] )
#define OFF(p)     ( ((unsigned *) &(p)) [0] )
```

Taking `SEG()` as characteristic, &(p) creates a temporary variable of type pointer-to-pointer which holds the address of that pointer in memory. &(p) can be treated as if it were the base address of an array of pointers—`*(&p)` and `(&p)[0]` both evaluate to p itself. The cast of (**unsigned** *) changes things, however. Now instead of being treated as an array of pointers, `&p` is treated as a two-element array of 16-bit **unsigned ints**, which are half the size of a pointer. The [1] picks up 16 bits of the 32-bit pointer, in this case the segment half. The [0] gets the offset half. These two macros are not at all portable, but they're more efficient than shift-and-mask operations, which are.

The `NBITS(type)` macro on line 68 of Listing A.1 evaluates to the number of bits in a variable of a particular integral type on the current machine. (`NBITS(int)` evaluates to the number of bits in an `int`.) It assumes that this number is 8, 16, 32, or 64, but the macro is easily modified to accept other values. The `_IS` macro on the previous line tests for one of these possibilities:

Number of bits required
for a type, NBITS().

```
#define IS(t,n) (((t)1 << (n))!=0)
```

It is passed the type (t) and the current guess (n). It casts the number 1 into the correct type, and then shifts that number to the left by the required number of bits. The resulting binary number for several shift values and a 16-bit-wide type are shown below:

n	$1 << n$
0	00000000000000000000 1
4	000000000000 10000
8	0000000100000000
12	0001000000000000
16	0000000000000000
32	0000000000000000

The expression evaluates to 0 if a variable of type `t` contains fewer than `n` bits. The `NBITS` macro, adds together several terms, one for each group of four bits in the number. For example, when an `int` is 16 bits wide, `_IS(int, 4)`, `_IS(int, 8)`, and `_IS(int, 12)` all evaluate to 1. The other terms evaluate to 0. Multiplying by four adjust the result to an actual bit count. I should point out that this macro, though huge, usually evaluates to a single constant at compile time. Since all the arguments are constants, none of the arithmetic needs to be done at run time. It's size might cause problems with your compiler's running out of macro-expansion space, however. Also, note that some compilers use greater precision for compile-time calculations than is available at run time—`NBITS()` won't work on these compilers.

`(unsigned) 0` evaluates to an `unsigned int` with all bits set to 1. The left shift

Finding the largest integer, MAXINT.

Array-manipulation macros:
NUMELE(),
LASTELE(), **TOOHIGH()**,
TOOLOW(), **INBOUNDS()**.

Maximum and minimum values, **max()**, **min()**,
RANGE().

Create a stack.

puts a zero into the high bit—the cast to (**unsigned**) defeats any potential sign extension on the right shift.

The definitions on lines 58 to 62 of Listing A.1 are for array manipulations. **NUMELE** is passed an array name and evaluates to the number of elements in the array. **LASTELE** evaluates to the address of the last element in the array. **TOOHIGH(array, p)** evaluates to true only if the pointer **p** is above the array. **TOOLOW(array, p)** is true if **p** is beneath the start of the array.² **INBOUNDS(array, p)** is true if **p** is in bounds.

The macros on lines 73 to 79 of *debug.h* test for various ranges of integer values. **max(a, b)** evaluates to **a** or **b**, whichever is larger; **min(a, b)** evaluates to **a** or **b**, whichever is smaller; and **RANGE(a, b, c)** evaluates to true if $a \leq b \leq c$.

A.1.2 *stack.h* and *yystack.h*—Generic Stack Maintenance

Compilers tend to use stacks a lot. To simplify stack maintenance, a series of macros that declare and manipulate stacks of an arbitrary, nonaggregate type, are declared in *stack.h*. The macros are used as follows:

`stack_dcl(stack, type, size)`

This macro creates a stack consisting of `size` objects of the given `type`. The `type` can be a simple type (like **int** and **long**) or a pointer (like **char***, **long****, and **struct foo ***). Stacks of aggregate objects like structures and arrays are not supported, however. Note that very complex types (like pointers to arrays) are best defined with a previous **typedef**, like this:

```
typedef int (*ptr_to_10_element_array)[10];
stack_dcl( stack, ptr_to_10_element_array, 128 );
```

otherwise they might not compile.

The `stack_dcl` macro declares three objects, the names of which are all derived from the `stack` argument as follows:

```
stack
p_stack
t_stack
```

If you put a definition of the form

```
#undef stack_cls
#define stack_cls static
```

somewhere after the **#include <tools/stack.h>**, then the `stack` and `stack` pointer will have the indicated storage class (**static** here).

2. Though ANSI requires that a pointer be able to point just past the end of an array, it does not allow a pointer to go past the start of the array. This behavior has serious practical ramifications in the 8086 compact and large models, and you won't be able to use **TOOLOW** with these memory models. Most 8086 compilers assume that no data object in the compact or large model is larger than 64K. Consequently, once the initial segment-register load is performed, only the offset portion of the address needs to be checked or modified when a pointer is modified. This practice causes problems when the base address of the array is close to the bottom of a segment. If, for example, an integer array is based at 2000:0000, a pointer increment correctly yields the address 2000:0002, but a decrement will yield 2000:FFFE rather than 1000:FFFE because the segment portion of the address won't be modified. This incorrect pointer is treated as very large, rather than very small.

`stack_p(stack)`

Access stack pointer directly.

This macro evaluates to the name used internally for the stack pointer. (See also: `stack_ele()`, below.)

`push(stack, x)`

Push an element.

Push `x` onto the stack. The `stack` argument must have appeared in a previous `stack_dcl()` invocation. This macro checks for stack overflow and invokes the `stack_err()` macro, discussed below, if the stack is full. (The default action prints an error message and terminates the program.) The `push_(stack, x)` does no overflow checking, so it is faster.

`pop(stack)`

Pop an element.

Pop an item off the stack (like this: `x = pop(stack);`). The `stack` argument must have appeared in a previous `stack_dcl()` invocation. This macro checks for stack underflow and invokes the `stack_err` macro, discussed below, if the stack is empty. (The default action prints an error message and terminates the program.) The `pop_(stack)` macro does no underflow checking, so it is faster.

`popn(stack, amt)`

Pop multiple elements.

This macro pops `amt` elements from the indicated stack, and evaluates to the element that was at the top of stack before anything was popped. There's also a `popn_(stack, amt)` macro, which does no underflow checking, but is otherwise the same as `popn()`.

`stack_clear(stack)`

Reinitialize stack.

This macro sets the stack back to its initial, empty condition, discarding all stack elements.

`stack_ele(stack)`

Get number of elements.

This macro evaluates to the number of elements currently on the indicated stack.

`stack_empty(stack)`

Test for stack empty.

This macro evaluates to true if the stack is empty.

`stack_full(stack)`

Test for stack full.

This macro evaluates to true if the stack is full.

`stack_item(stack, offset)`

Access arbitrary element.

Evaluates to the item at the indicated offset from the top of the stack. Note that no range checking is done to see if `offset` is valid. Use `stack_ele()` to do this checking if you need it.

`stack_err(overflow)`

Stack error handler.

This macro is invoked to handle stack errors. The `overflow` argument is 1 if a stack overflow occurred (you tried to push an item onto a stack that was full). It's zero on an underflow (you tried to pop an item off an empty stack). The default action prints an error message and terminates the program, but you can `#undef` this macro and redefine it to do something more appropriate if you like. If this macro evaluates to something, that value will, in turn, come back from the `push` or `pop` macros when an error occurs. For example:

```
#undef stack_err
#define stack_err(overflow) 0
```

will cause both push and pop to evaluate to 0 if an overflow or underflow occurs.

Stack-macro implementation.

The ANSI concatenation operator, ##.

Downward-growing stacks are used for ANSI compatibility.

Pushing and popping, push(), push_(), pop(), pop_(), popn(), popn_().

Stack-macro error processing, stack_err().

Double cast needed to convert int to pointer.

Stack macros for `oecs`, `<yyystack.h>`.

The stack macros are implemented in *stack.h*, Listing A.2. These macros make heavy use of the ANSI concatenation operator, ##. This operator is removed from a macro definition by the preprocessor, effectively concatenating the strings on either side of it. For example, given the invocation

```
stack_dcl(plates, int, 128);
```

the `typedef` on line seven of Listing A.2 evaluates, first to

```
typedef int t_##plates;
```

(the `int` replaces `type` in the definition, the `plates` replaces `stack`) and then the ## is removed, yielding:

```
typedef int t_plates;
```

Lines eight and nine will evaluate as follows:

```
t_plates plates[128];
t_plates *(p_plates) = plates + (128);
```

given the earlier `stack_dcl()` invocation. A downward-growing stack is created, so the stack pointer (`p_plates`) is initialized to just past the end of the stack. This initialization works correctly in ANSI C, which says that pointer can point one cell beyond the end of an array and still be valid, but a pointer is not permitted to go to the left of the array.

Since a downward-growing stack is implemented, a push is a predecrement (on line 23) and a pop is a postincrement (on line 31). This way the stack pointer always points at the top-of-stack item. The actual push and pop operations have been isolated into the `push_()` and `pop_()` macros because it's occasionally useful to do a push or pop without checking for stack overflow or underflow. Note that the `popn_()` macro on line 34 must evaluate to the top-of-stack element before anything is popped. A postincrement is simulated here by modifying the stack pointer by `amt` elements, and then reaching backwards with a `[-amt]` to get the previous top-of-stack element.

The `stack_err()` macro on lines 39 and 40 use `ferr` to print an error message and then exits the program. (`ferr()` works like `printf()` except that it sends characters to `stderr` and calls `exit(1)` when it's done. It is discussed in depth below). In theory, the entire expression evaluates to `ferr()`'s return value, but since `ferr()` doesn't return, it's not really an issue here.

You need two casts on lines 27 and 31 where `stack_err()` is invoked because there's no telling what sort of value to which the macro will evaluate. The default `ferr()` call evaluates to an `int`, but since `stack_err()` can be redefined, you can't count on that. The stack type is also uncertain because it can be redefined. You can't just cast `stack_err()` to the same type as the stack because the stack might be an array of pointers—many compilers complain if you try to convert an `int` to a pointer. So you have to cast twice—first to `long` and then to the stack type—to suppress this error message. Though casting to `long` as an intermediate step fixes this problem, though it might introduce an unnecessary conversion if the stack type is an `int`, `short`, or `char`. These multiple type conversions will also cause portability problems if the `stack_err()` macro evaluates to something that won't fit into a `long` (like a `double`).

A second set of stack macros are in the file `<yyystack.h>`, Listing A.3. These macros differ from the ones just described in three ways only: all the macro names are preceded by the characters `yy`, the generated names are prefaced by `yyt_` and `yyp_` rather

Listing A.2. *stack.h*—Stack-Maintenance Macros

```

1  /* Stack.h      Stack-maintenance macros. Creates downward-growing stacks
2   *          (which should work in all six memory models).
3   */
4
5 #define stack_cls    /* empty */
6
7 #define stack_dcl(stack,type,size)  typedef type t_##stack;
8                                stack_cls t_##stack stack[size];
9                                stack_cls t_##stack (*p_##stack)
10                               = stack + (size)
11 #define stack_clear(stack) ( (p_##stack) = (stack + \
12                           sizeof(stack)/sizeof(*stack)) )
13
14 #define stack_full(stack) ( (p_##stack) <= stack )
15 #define stack_empty(stack) ( (p_##stack) >= (stack + \
16                           sizeof(stack)/sizeof(*stack)) )
17
18 #define stack_ele(stack) ((sizeof(stack)/sizeof(*stack)) - (p_##stack-stack))
19
20 #define stack_item(stack,offset) ( *(p_##stack + (offset)) )
21 #define stack_p(stack)           p_##stack
22
23 #define push_(stack,x) ( **p_##stack = (x) )
24 #define pop_(stack)   ( *p_##stack++)
25
26 #define push(stack,x) ( stack_full(stack)
27             ? ((t_##stack) (long) (stack_err(1)))
28             : push_(stack,x)
29
30 #define pop(stack)   ( stack_empty(stack)
31             ? ((t_##stack) (long) (stack_err(0)))
32             : pop_(stack)
33
34 #define popn_(stack,amt) ( (p_##stack += amt)[-amt] )
35 #define popn(stack,amt) ( (stack_ele(stack) < amt)
36             ? ((t_##stack) (long) (stack_err(0)))
37             : popn_(stack,amt)
38
39 #define stack_err(o) ((o) ? ferr("Stack overflow\n") \
40                         : ferr("Stack underflow\n"))

```

than *t_* and *p_*, and I've abbreviated *yystack* to *yystk*. The second set of macros are used only within the **LLama**- and **occs**-generated parsers. The extra *yy* avoids potential name conflicts with user-supplied names.

A.1.3 *l.h* and *compiler.h*

Two other include files are of interest: *l.h* contains prototypes for the routines in *l.lib*, (which contains run-time subroutines for **LEX** and **occs**-generated programs); *compiler.h* contains prototypes for the routines in *comp.lib* (which contains subroutines used by **LEX** and **occs** themselves). These files are not listed here because they are created automatically from within the makefile that creates the library. The Microsoft compiler's /Zg switch—which outputs prototypes for all functions in the input file rather than compiling the file—is used. Copies of the prototype files are included on the distribution disk, however, and they are **#included** by most of the source files in this book.

Listing A.3. *yystack.h*— Stack-Maintenance Macros: LLama/Occs Version

```

1  /* yystack.h      Stack-maintenance macros --- yacc and llama version.
2   */
3
4  #define yystk_cls      /* empty */
5
6  #define yystk_dcl(stack,type,size)  typedef type yyt_##stack;           \
7                                yystk_cls yyt_##stack stack[size];          \
8                                yystk_cls yyt_##stack (*yyp_##stack)    \
9                                = stack + (size)
10 #define yystk_clear(stack) ( (yyp_##stack) = (stack + \
11                           sizeof(stack)/sizeof(*stack)) )
12
13 #define yystk_full(stack) ( (yyp_##stack) <= stack )
14 #define yystk_empty(stack) ( (yyp_##stack) >= (stack + \
15                           sizeof(stack)/sizeof(*stack)) )
16
17 #define yystk_ele(stack)  ((sizeof(stack)/sizeof(*stack)) - \
18                           (yyp_##stack-stack))
19
20 #define yystk_item(stack,offset) ( *(yyp_##stack + (offset)) )
21 #define yystk_p(stack)        yyp_##stack
22
23 #define yypush_(stack,x) ( *--yyp_##stack = (x) )
24 #define yypop_(stack)        (*yyp_##stack++)
25
26 #define yypush(stack,x)     ( yystk_full(stack) \
27                           ? ((yyt_##stack)(long) (yystk_err(1))) \
28                           : yypush_(stack,x) )
29
30 #define yypop(stack)       ( yystk_empty(stack) \
31                           ? ((yyt_##stack)(long) (yystk_err(0))) \
32                           : yypop_(stack) )
33
34 #define yypopn_(stack,amt) ( (yyp_##stack += amt)[-amt] )
35 #define yypopn(stack,amt) ( (yystk_ele(stack) < amt) \
36                           ? ((yyt_##stack)(long) (yystk_err(0))) \
37                           : yypopn_(stack,amt) )
38
39 #define yystk_err(o)        ((o) ? ferr("Stack overflow\n") \
40                           : ferr("Stack underflow\n"))

```

A.2 Set Manipulation

Bit maps.

Many of the operations involved in compiler writing—like creating state-machine tables from regular expressions and creating bottom-up parse tables from a grammar—involve operations on sets, and C, unlike Pascal, doesn't have a built-in set capability. Fortunately, it's not too hard to implement sets in C by means of bit maps—one-dimensional arrays of one-bit numbers. This section presents a package of bit-map-based, set-manipulation routines.

A.2.1 Using the Set Functions and Macros

`<tools/set.h>`, SET.

To use the set routines, you must put a **#include** `<tools/set.h>` at the top of your file. The definitions for a SET are found there, and many of the set functions are actually macros defined in `<tools/set.h>`. The set routines are described in the following paragraphs. I'll give a few examples after the calling-conventions are described. If a

name is in all caps in the following list, it's implemented as a macro, and most of these macros have side effects. Be careful.

SET *newset (**void**)

Create a set.

Create a new set and return a pointer to it. Print an error message and raise SIGABRT if there's insufficient memory. Normally this signal terminates the program but you can use `signal()` to change the default action (the process is described in a moment). NULL is returned if `raise()` returns.

void delset (SET *set)

Delete a set.

Delete a set created with a previous `newset` call and free the associated memory. The argument must have been returned from a previous `newset ()` call.

SET *dupset (SET *set)

Duplicate a set.

Create a new set that has the same members as the input set. This routine is more efficient than using `newset ()` to create the set and then copying the members one at a time, but otherwise has the same effect.

int _addset (SET *set, **int** bit)

Internal set function: enlarge set.

This is an internal function used by the ADD macro, and shouldn't be called by you.

int num_ele (SET *set)

Find number of elements in set.

Return the number of elements in the input set. NULL sets (described below) are considered to be empty.

int _set_test (SET *set1, SET *set2)

Internal set function: test bits.

This is another workhorse function used internally by the macros. Don't call it directly.

int setcmp (SET *set1, SET *set2)

Compare two sets.

Compare two sets in a manner similar to `strcmp()` returning 0 if the sets are equivalent, <0 if `set1`<`set2` and >0 if `set1`>`set2`. This routine lets you sort an array of SET's so that equivalent ones are adjacent. The determination of less than and greater than is pretty much arbitrary. (The routine just compares the bit maps as if you were doing a lexicographic ordering of an array of ints.)

unsigned sethash (SET *set1)

Compute hash value for set.

This function is even more obscure than `setcmp()`. It is provided for those situations where a SET is used as the key in a hash table. It returns the sum of the individual words in the bit map.

int subset (SET *set, SET *sub)

Sub is subset of set.

Return 1 if `sub` is a subset of `set`, 0 otherwise. Empty and null sets are subsets of everything, and 1 is returned if both sets are empty or null.

void _set_op (**int** op, SET *dest, SET *src)

Internal set function: binary operations.

Another workhorse function used internally by the macros.

void invert (SET *set)

Complement set by inverting bits.

Physically invert the bits in the set, setting 1's to 0 and vice versa. In effect, this operation removes all existing members from a set and adds all possible members

Clear set and make smaller.

void truncate(SET *set)

Clears the set and shrinks it back to the original, default size. Compare this routine to the CLEAR() macro, described below, which clears all the bits in the map but doesn't modify the size. This routine is really a more efficient replacement for delset(s); s=newset();. If the original set isn't very big, you're better off using CLEAR().

Find next set element.

int next_member(SET *set)

When called several successive times with the same argument, returns the next element of the set each time it's called or -1 if there are no more elements. Every time the set argument changes, the search for elements starts back at the beginning of the set. A NULL argument also resets the search to the beginning of the set (and does nothing else). Strange things happen if you add members to the set between successive calls. If calls to next_member() are interspersed with calls to pset() (discussed below), next_member() won't work properly. Calls to next_member() on different sets cannot be interspersed.

Print set elements.

void pset(SET *set, **int** (*out)(), **void** *param)

Print the set. The output routine pointed to by out is called for each element of the set with the following arguments:

(*out)(param, "null", -1); Null set
 (*out)(param, "empty", -2); Empty set
 (*out)(param, "%d ", N); Set element N

This way you can use fprintf() as a default output routine.

dest = dest \cup src.

UNION(SET *dest, SET *src)

Modify the dest set to hold the union of the src and dest sets.

dest = dest \cap src.

INTERSECT(SET *dest, SET *src)

Modify the dest set to hold the intersection of the src and dest sets.

dest = symmetric difference.

DIFFERENCE(SET *dest, SET *src)

Modify the dest set to hold the symmetric difference of the src and dest sets. (An element is put into the target set if it is a member of dest but not of src, or vice versa.)

dest = src.

ASSIGN(SET *dest, SET *src)

Overwrite the dest with src.

Clear set.

CLEAR(SET *s)

Clear all bits in s, creating an empty set.

Add all elements to set.

FILL(SET *s)

Set all bits in s to 1, creating a set that holds every element in the input alphabet.

Complement set.

COMPLEMENT(SET *s)

Complement a set efficiently by modifying the set's complement bit. Sets

complemented in this way can not be manipulated by UNION(), etc. See also, invert() and INVERT().

INVERT(SET *s)

Invert all bits in bit map.

Complement a set by physically changing the bit map (see text).

IS_DISJOINT(SET *s1, SET *s2)

Test for disjoint.

Evaluate to true only if the two sets are disjoint (have no elements in common).

IS_INTERSECTING (SET *s1, SET *s2) Test for intersection.

Evaluate to true only if the two sets intersect (have at least one element in common).

IS_EMPTY(SET *s)

Test for empty set.

Evaluate to true only if set is empty (having no elements) or null (s is NULL)

IS_EQUIVALENT(SET *s1, SET *s2)

Text for equivalence.

Evaluate to true only if the two sets are equivalent (have the same elements).

ADD(SET *s, int x)

Add member to set.

Add the element c to set s. It is not an error to add an element to a set more than once.

REMOVE(SET *s, int x)

Remove member from set.

Remove the element c from set s. It is not an error to remove an element that is not in the set.

TEST(SET *s, int x)

Test for membership (all sets).

Evaluates to true if x is an element of set s.

MEMBER(SET *s, int x)

Test for membership (no complemented sets).

Evaluates to true if x is an element of set s. This macro doesn't work on COMPLEMENTED sets, but it's both faster and smaller than TEST, which does. The distinction is described below.

The elements of sets must be numbers, though in many instances any arbitrary number will do. Enumerated types are almost ideal for this purpose, though #defines can be used too. For example:

```
typedef enum
{
    JAN, FEB, MAR,
    APR, MAY, JUN,
    JUL, AUG, SEP,
    OCT, NOV, DEC
} MONTHS;
```

creates 12 potential set elements. You can create two sets called winter and spring by using the following set operations:

Sets, an example. Set elements are numbers.

```
#include <set.h>

SET *winter, *spring;

winter = newset();
spring = newset();

ADD( JAN, winter );
ADD( FEB, winter );
ADD( MAR, winter );
ADD( APR, spring );
ADD( MAY, spring );
ADD( JUN, spring );
```

Set operations can now be performed using the other macros in `<tools/set.h>`. For example: `IS_DISJOINT(winter, spring)` evaluates to true because the sets have no elements in common; `IS_EQUIVALENT(winter, spring)` evaluates to false for the same reason. A third set that contains the union of `spring` and `winter` can be created with:

```
half_year = dupset( winter );
UNION( half_year, spring );
```

Something like:

```
half_year = dupset( winter );
INTERSECT(half_year, spring);
```

creates an empty set because there are no common elements.

Set implementation difficulties. Null and empty sets.

There are two implementation difficulties with the set routines. The first is the difference between a null set and an empty set. (I'll bet that you thought that the difference was just one more obscure mathematical conundrum designed for no other purpose than to make undergraduates' heads swim). An empty set is a set that has no elements. In the case of the routines presented here, `newset()` and `dupset()` both create empty sets. They have allocated an internal data structure for representing the set, but that set doesn't have anything in it yet. A null set, however, is a `SET` pointer with nothing in it. For example:

```
SET *p = NULL; /* p represents the null set */
p = newset(); /* p now represents an empty set */
```

Problems with complemented sets.

In practice, this difference means that the routines have to be a bit more careful with pointers than they would be otherwise, and are a little slower as a consequence.

Complimented sets present another problem. You'll notice that the eventual size of the set doesn't have to be known when the set is created by `newset()`. The set size is just expanded as elements are added to it. This can cause problems when you complement a set, because the complemented set should contain all possible elements of input alphabet except the ones that are in the equivalent, uncomplemented, set. For example, if you're working with a language that's comprised of the set of symbols {A, B, C, D, E, F, G} and you create a second set {A,C,E,G} from elements of the language, the complement of this second set should be {B,D,F}.

Complement by physically inverting bits.

This ideal situation is difficult to do, however. Sets are represented internally as bit maps, and these maps are of finite size. Moreover, the actual size of the map grows as elements are added to the set. You can complement a set by inverting the sense of all the bits in the map, but then you can't expand the set's size dynamically—at least not without a lot of work. To guarantee that a complemented set contains all the potential elements, you first must expand the set size by adding an element that is one larger than any possible legitimate element, and then complement the expanded set. A second

problem has to do with extra elements. The bit-map size is usually a little larger than the number of potential elements in the set, so you will effectively add members to the set if you just stupidly complement bits. On the plus side, set operations (union, intersection, etc.) are much easier to do on physically complemented sets.

An alternate method of complementing the set uses negative-true sets and positive-true sets. Here, you mark a set as negative or positive by setting a bit in the SET structure. You don't have to modify the bit map at all. If a set is marked negative-true when you test for membership, you can just reverse the sense of the test (evaluate to true if the requested bit is not false). Though this method solves the size problem, operations on negative-true sets are much harder to perform.

Complement by marking
set as negative true.

Since the two representations are both useful, but in different applications, I've implemented both methods. The `INVERT()` macro performs a ones-complement on all bits currently in the bit map. Note that if new elements are added, the new bits won't be complemented. You should always expand a set out to the maximum number of elements (by adding and then removing the largest element) before inverting it. The `COMPLEMENT()` macro implements the second method. It doesn't modify the bit map at all; rather, it sets a flag in the `SET` structure to mark a set as negative true.

Physical complement,
INVERT()

Because there are two different classes of sets (those that are complemented and those that are inverted), there are also two different macros for testing membership. MEMBER() evaluates to true only if the bit corresponding to the requested element is actually set to one. MEMBER() can't be used reliably on complemented sets. The TEST() macro can be used with complemented sets, but it's both larger and slower than MEMBER(). If a set is complemented, the sense of the individual bits is reversed as part of the testing process. If the set isn't complemented TEST() works just like MEMBER().

Logical complement, COMPLEMENT ()

Testing for membership
on complemented sets,
`MEMBER()` vs. `TEST()`.

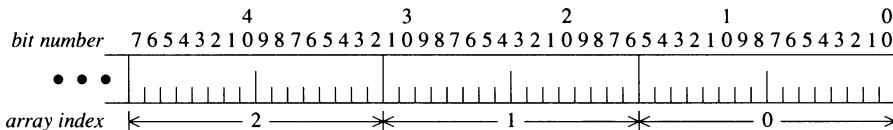
The various set operations (UNION, INTERSECT, and so forth) are only valid on inverted sets. Use INVERT() if you're going to perform subsequent operations on the inverted set. I leave it as an exercise to the reader to make `_set_op()` work on complemented sets. The complement bit can represent all bits not in the bit map—if the complement bit is set, all bits not in the map are zero and vice versa. I've found these routines quite workable in their existing state—it seemed pointless to make the code both larger and slower to correct what has turned out not to be a problem.

Set operations work only on physically implemented sets.

A.2.2 Set Implementation

The set routines are implemented in two places, the first of which is the macro file `<tools/set.h>`, in Listing A.4. The file starts with various system-dependent definitions, concentrated in one place to make it easier to port the code. `_SETTYPE` (on line three) is used as the basic unit in a bit map. As I mentioned earlier, a bit map is a 1-dimensional array of 1-bit objects. In terms of set operations, a number is in the set if the array element at that position in the map is true (if `bitmap[5]` is true than 5 is in the set). The map is implemented using an array of `SETTYPES` as follows:

<tools/set.h>, SETTYPE.



Implementing bit maps.

Bit 0 is at position 0 in array[0], bit 1 is at position 1 in array[0], bit 20 is at position 4 in array[1], and so forth. In order to make the array manipulation as efficient as possible, `_SETTYPE` should be the largest integral type that can be manipulated with a single instruction on the target machine. In an 8086, for example, the largest such type is a 16-bit word—larger numbers require several instructions to manipulate them—so I've used the 16-bit `short` as the `_SETTYPE`.

Listing A.4. *set.h*—Macro Definitions and Prototypes for Set Functions

```

1  /* SET.H:  Macros and function prototypes for the set functions */
2
3  typedef unsigned short      _SETTYPE;           /* one cell in bit map */
4  #define _BITS_IN_WORD        16
5  #define _BYTES_IN_ARRAY(x)   (x << 1)          /* # of bytes in bit map */
6  #define _DIV_WSIZE(x)        ((unsigned) (x) >> 4)
7  #define _MOD_WSIZE(x)        ((x) & 0x0f)
8  #define _DEFWORDS 8          /* elements in default set */
9  #define _DEFBITS  (_DEFWORDS * _BITS_IN_WORD) /* bits in default set */
10 #define _ROUND(bit)  (((_DIV_WSIZE(bit) + 8) >>3 ) <<3 )
11
12 typedef struct _set_
13 {
14     unsigned char nwords;           /* Number of words in map */
15     unsigned char compl;           /* is a negative true set if true */
16     unsigned nbits;              /* Number of bits in map */
17     _SETTYPE *map;               /* Pointer to the map */
18     _SETTYPE defmap[ _DEFWORDS ]; /* The map itself */
19
20 } SET;
21
22 extern int      _addset;       P(( SET*, int
23 extern void    delset;        P(( SET*)
24 extern SET    *dupset;       P(( SET*)
25 extern void    invert;        P(( SET*)
26 extern SET    *newset;       P(( void
27 extern int      next_member;  P(( SET *)
28 extern int      num_ele;      P(( SET*)
29 extern void    pset;         P(( SET*, int (*)(), void*
30 extern void    _set_op;       P(( int,           SET*, SET*)
31 extern int      _set_test;    P(( SET*, SET*)
32 extern int      setcmp;       P(( SET*, SET*)
33 extern unsigned sethash;      P(( SET*)
34 extern int      subset;       P(( SET*, SET*)
35 extern void    truncate;     P(( SET*)
36
37                                     /* Op argument passed to _set_op */
38 #define _UNION          0      /* x is in s1 or s2 */
39 #define _INTERSECT       1      /* x is in s1 and s2 */
40 #define _DIFFERENCE     2      /* (x in s1) and (x not in s2) */
41 #define _ASSIGN          4      /* s1 = s2 */
42
43 #define UNION(d,s)      _set_op( _UNION,      d, s )
44 #define INTERSECT(d,s)  _set_op( _INTERSECT,  d, s )
45 #define DIFFERENCE(d,s) _set_op( _DIFFERENCE, d, s )
46 #define ASSIGN(d,s)     _set_op( _ASSIGN,     d, s )
47
48 #define CLEAR(s)         memset( (s)->map, 0, (s)->nwords * sizeof(_SETTYPE) )
49 #define FILL(s)          memset( (s)->map, ~0, (s)->nwords * sizeof(_SETTYPE) )
50 #define COMPLEMENT(s)   ( (s)->compl = ~ (s)->compl )
51 #define INVERT(s)        invert(s)
52
53 #define _SET_EQUIV        0      /* Value returned from _set_test, equivalent */
54 #define _SET_DISJ         1      /* disjoint */
55 #define _SET_INTER        2      /* intersecting */
56
57 #define IS_DISJOINT(s1,s2) ( _set_test(s1,s2) == _SET_DISJ )
58 #define IS_INTERSECTING(s1,s2) ( _set_test(s1,s2) == _SET_INTER )
59 #define IS_EQUIVALENT(a,b) ( setcmp((a),(b)) == 0 )

```



Listing A.4. continued...

```

60 #define IS_EMPTY(s)           ( num_ele(s) == 0 )
61
62 /* All of the following have heavy-duty side-effects. Be careful. */
63
64 #define _GBIT(s,x,op) ( ((s)->map)[_DIV_WSIZE(x)] op (1 << _MOD_WSIZE(x)) )
65
66 #define REMOVE(s,x)   (((x) >= (s)->nbits) ? 0 : _GBIT(s,x,&= ~) )
67 #define ADD(s,x)      (((x) >= (s)->nbits) ? _addset(s,x) : _GBIT(s,x,|= ) )
68 #define MEMBER(s,x)   (((x) >= (s)->nbits) ? 0 : _GBIT(s,x,& ) )
69 #define TEST(s,x)     (( MEMBER(s,x) ) ? !(s)->compl : (s)->compl )

```

The macros on lines four to seven all reflect the size of the `_SETTYPE` and will have to be changed if `_SETTYPE` is changed. `_BITS_IN_WORD` is just that, the number of bits in a variable of type `_SETTYPE`. `_BYTES_IN_ARRAY` is passed a count, and it returns the number of bytes in an array of that many `_SETTYPE`-sized variables. In this case, I multiply the count by two using a left shift rather than doing something like `s * sizeof(_SETTYPE)`, which uses a multiply on my compiler — the shift is a more efficient operation than the multiply.

The `_DIV_WSIZE(x)` and `_MOD_WSIZE(x)` macros help determine the position of a particular bit in the map. The argument to both macros is the bit for which you want the position. The `_DIV_WSIZE(x)` macro evaluate to the array element that holds the bit. A `_DIV_WSIZE(20)` evaluates to 1 because bit 20 is in array[1]. It's just doing an integer divide by 16, though I'm using a left shift here for efficiency's sake. The `_MOD_WSIZE(x)` macro evaluate to the position of the bit within the word—the offset in bits from the least-significant bit in the word. A `_MOD_WSIZE(20)` evaluates to 4 because bit 20 is at offset 4 of array[1]. I'm doing an efficient modulus-16 operation by using a bitwise AND rather than a %.

`_DEFWORDS`, on line eight of Listing A.4, determines the array size for a default bit map. Initially, all bit maps have this many elements. The size is expanded (in `_DEFWORDS`-sized chunks) if needed. `_DEFWORDS` is set to 8 here, so the default map can have 128 elements before it needs to be expanded. `_DEFBITS`, on the next line, is just the number of bits required for that many words.

The `_ROUND` macro is used to expand the size of the array. The array grows in `_DEFWORDS`-sized chunks. Say, for example, that the array starts out at the default size of 8 words, and you want to add the number 200 to the set. The array must be expanded to do so, and after the expansion, the array should have 16 elements in it ($2 \times _DEFWORDS$). In this situation the macro expands to:

```
(((_DIV_WSIZE(200) + 8) >>3) <<3)
```

and one more level to:

```
((((unsigned)(200) >> 4) + 8) >>3) <<3)
```

The $200 >> 4$ evaluates to 12, so bit 200 is in array[12]. 12 plus 8 is 20, and $20 >> 3$ yields 2. (The $>> 3$ is an integer divide by 8.) The final multiply by 8 (the $<< 3$) yields 16, so the map array is expanded to 16 elements.

The SET itself is represented with the following structure, defined on lines 12 to 20 of Listing A.4.

`_BITS_IN_WORD`,
`_BYTES_IN_ARRAY`.

Computing bit positions,
`_DIV_WSIZE(x)`,
`_MOD_WSIZE(x)`.

`_DEFWORDS`, `_DEFBITS`,
`_ROUND`.

Representing a set, SET.

```
typedef struct _set_
{
    unsigned char nwords ; /* Number of words in map */
    unsigned char compl ; /* negative true if set */
    unsigned nbits ; /* Number of bits in map */
    _SETTYPE *map ; /* Pointer to the map */
    _SETTYPE defmap[_DEFWORDS] ; /* The default map */
}

} SET;
```

The number of bits in the map (nbts) could be computed from the number of words (nwords), but it's more efficient—of computation time, not space—to keep both numbers in the structure. The compl field is set to true if this is a negative-true set. The defmap[] array is the default bit map. Initially, map just points to it. When the map grows, however, a new array is allocated with malloc() and map is modified to point at the new array. This strategy is used rather than a realloc() call for run-time efficiency, at the cost of wasted memory—realloc() will generally have to copy the entire structure, but only the map needs to be copied if you do it yourself.

Set macros.

CLEAR, FILL.

Adding and removing set members, ADD(), REMOVE(). Testing for membership, MEMBER().

_GBIT.

<set.h> continues on line 22 of Listing A.4 with function prototypes for the actual functions. The macros on lines 43 to 60 handle operations that actually modify a set (union, intersection, symmetric difference, and so forth). Most of them map to set_op() calls, passing in a constant to tell the subroutine which function is required. CLEAR and FILL modify the bit map directly, however, calling memset() to fill the bits with zeros or ones as appropriate.

The real heart of <set.h> are the set-manipulation macros at the end of the file, reproduced here:

```
#define _GBT(s,x,op) (((s)->map)[_DIV_WSIZE(x)] op (1 << _MOD_WSIZE(x)))

#define REMOVE(s,x) (((x) >= (s)->nbts) ? 0 : _GBT(s,x,&= ~))
#define ADD(s,x) ((((x) >= (s)->nbts) ? _addset(s,x) : _GBT(s,x,|= )))
#define MEMBER(s,x) (((x) >= (s)->nbts) ? 0 : _GBT(s,x,& ))
#define TEST(s,x) (( MEMBER(s,x) ) ? !(s)->compl : (s)->compl )
```

The REMOVE, ADD, and MEMBER macros all evaluate to _GBT invocations. The only difference is the operator (op) passed into the macro. The first part of _GBT uses the following to select the array element in which the required bit is found:

```
( (s)->map )[ _DIV_WSIZE(x) ]
```

The second half of the macro shifts the number 1 to the left so that it will fall in the same position as the required bit, using

```
(1 << _MOD_WSIZE(x)) )
```

If you're accessing bit 5, the number 1 is shifted left 5 bits, yielding the following binary mask in a 16-bit word:

```
0000000000100000
```

The same shift happens when you access bit 20, but in this case the first half of the macro chooses (s)->map[1] rather than (s)->map[0]. The op argument now comes into play. The shifted mask is ORed with the existing array element to add it to the set. Similarly, the inverse of the mask is ANDed to the array element to clear a bit. OR is also used to test for membership, but since there's no = in MEMBER, the map is not modified.

The TEST macro, which can handle complemented sets, works by first determining whether x is in the set, and then evaluating to the complement flag or its inverse, as appropriate. For example, the complemented flag will be 1 if x is in a negative true set,

Testing complemented sets for membership, TEST().

and MEMBER tests true if the bit is in the map. The inverse of the complement flag (0) is returned, however. I've used this somewhat convoluted approach because of the sheer size of the macro. TEST(s, x) expands to this monster:

```
(((( (x) >= (s)->nbits) ? 0 :
    (((s)->map)[((unsigned)(x)>>4)] & (1<<((x) & 0x0f)))) ?  
    !(s)->compl : (s)->compl);
```

but the more obvious solution:

```
(s)->compl ? !MEMBER(s,x) : MEMBER(s,x)
```

turns into this:

```
(s)->compl  
? !( ((x) >= (s)->nbits) ? 0 :  
    (((s)->map)[((unsigned)(x)>>4)] & (1<<((x) & 0x0f))))  
: ((x) >= (s)->nbits) ? 0 :  
    (((s)->map)[((unsigned)(x)>>4)] & (1<<((x) & 0x0f))))
```

which is even worse.

The functions needed for set manipulation are all in *set.c*, the first part of which is Listing A.5. The newset() function, which actually creates the new set, starts on line nine. Normally, a pointer to the newly allocated SET is returned. If insufficient memory is available, an error message is printed and the program is terminated by the raise(SIGABRT) call on line 21. raise() is an ANSI function which, in this case, causes the program to terminate. I've chosen to abort the program rather than return NULL (like malloc() does) because most applications will terminate the program anyway if memory isn't available. If you want NULL to be returned on an error rather than aborting the program, just disable the SIGABRT signal as follows:

Creating sets, newset().

raise().

```
#include <signal.h>
main()
{
    signal( SIGABRT, SIG_IGN );
}
```

Listing A.5. *set.c*—SET Creation and Destruction

```
1 #include <stdio.h>
2 #include <ctype.h>
3 #include <signal.h>
4 #include <stdlib.h>
5 #include <string.h>
6 #include <tools/debug.h>
7 #include <tools/set.h>
8
9 PUBLIC SET      *newset()
10 {
11     /* Create a new set and return a pointer to it. Print an error message
12      * and raise SIGABRT if there's insufficient memory. NULL is returned
13      * if raise() returns.
14      */
15
16     SET *p;
17
18     if( !(p = (SET *) malloc( sizeof(SET) )) )
19     {
20         fprintf(stderr,"Can't get memory to create set\n");
21         raise( SIGABRT );
22         return NULL;           /* Usually won't get here */
23 }
```



Listing A.5. continued...

```

23     }
24     memset( p, 0, sizeof(SET) );
25     p->map      = p->defmap;
26     p->nwords   = _DEFWORDS;
27     p->nbits    = _DEFBITS;
28     return p;
29 }
30
31 /*-----*/
32
33 PUBLIC void      delset( set )
34 SET        *set;
35 {
36     /* Delete a set created with a previous newset() call. */
37
38     if( set->map != set->defmap )
39         free( set->map );
40     free( set );
41 }
42
43 /*-----*/
44
45 PUBLIC SET      *dupset( set )
46 SET        *set;
47 {
48     /* Create a new set that has the same members as the input set */
49
50     SET      *new;
51
52     if( !(new = (SET *) malloc( sizeof(SET) )) )
53     {
54         fprintf(stderr,"Can't get memory to duplicate set\n");
55         exit(1);
56     }
57
58     memset( new, 0, sizeof(SET) );
59     new->compl  = set->compl;
60     new->nwords = set->nwords;
61     new->nbits  = set->nbits;
62
63     if( set->map == set->defmap )           /* default bit map in use */
64     {
65         new->map = new->defmap;
66         memcpy( new->defmap, set->defmap, _DEFWORDS * sizeof(_SETTYPE) );
67     }
68     else                                /* bit map has been enlarged */
69     {
70         new->map = (_SETTYPE *) malloc( set->nwords * sizeof(_SETTYPE) );
71         if( !new->map )
72         {
73             fprintf(stderr,"Can't get memory to duplicate set bit map\n");
74             exit(1);
75         }
76         memcpy( new->map, set->map, set->nwords * sizeof(_SETTYPE) );
77     }
78     return new;
79 }
```

The `del_set()` function starting on line 33 is the SET destructor subroutine. It frees any memory used for an expanded bit map on line 38, and then frees the memory used for the SET itself.

`dupset()` (starting on line 45 of Listing A.5) duplicates an existing set. It's much more efficient than calling `new_set()` and then adding members to the new set one at a time.

The functions in Listing A.6 handle set enlargement: `_addset()` is called from the ADD macro when the bit map is not large enough to hold the requested bit. All it does is call `enlarge()` (which starts on line 95) to make the map larger, and then invokes `_GBIT` to set the bit. `_enlarge()` is passed the required number of words in the bit map (`need`). The test on line 106 causes the routine to return if no expansion is required. Note that `exit()` is called on line 114 if `enlarge()` can't get memory. I've done this rather than call `raise()` because I'm assuming that the value returned from ADD will not be tested. It would be risky to call `raise()` in this situation because the signal handler might have been reassigned to an empty function, as was discussed earlier.

The next part of `set.c` (in Listing A.7) consists of various testing functions. The `num_ele()` function, starting on line 126, determines the number of elements in the set. It does this by looking at the map array one byte at a time, using a table lookup to do the counting. The table in question is `nbits[]`, declared on lines 134 to 152. The table is indexed with a number in the range 0–255, and it evaluates to the number of bits set to 1 in that number. For example, the decimal number 93 is 01011101 in binary. This number has five ones in it, so `nbits[93]` holds 5. The loop on lines 162 and 163 just goes through the map array byte by byte, looking up each byte in `nbits[]` and accumulating the count in `count`.

The `_set_test()` function starting on line 170 of Listing A.7 compares two sets, returning `_SET_EQUIV` if the sets are equivalent (have the same elements), `_SET_INTER` if the sets intersect (have at least one element in common) but aren't equivalent, and `_SET_DISJ` if the sets are disjoint (have no elements in common). A bitwise AND is used to test for intersection on line 203—the test for equivalence was done on line 195 of Listing A.7. If the AND tests true, there must be at least one bit at the same relative position in both bit-map elements.

Note that the sets are made the same size with the `enlarge()` calls on lines 187 and 188 before they are compared, and this expansion can waste time if the sets are likely to be different sizes. For this reason, a second comparison function, `setcmp()`, is provided on line 215. This routine works like `strcmp()`, returning zero if the sets are equivalent, a negative number if `set1 < set2`, and a positive number if `set1 > set2`. Since `setcmp()` does not modify the set sizes, using it can be less time consuming than using `_set_test()`. The main purpose of this second comparison function is to let you sort an array of SET pointers so that the equivalent sets are adjacent. The determination of less than and greater than is pretty much arbitrary—the routine just compares the maps as if they were arrays of `ints`. `setcmp()` first compares the map elements that exist in both of the sets on lines 228 to 230. If there's a mismatch, the two bytes are subtracted to get a relative ordering. The code on lines 236 to 247 is executed only if the two sets are identical up to this point. The tail end of the larger set is then scanned to make sure that it's all zeros. If so, the two sets are equivalent. If not, the larger set is the greater.

The final comparison function is `subset()`, starting on line 275 of Listing A.7. `common` and `tail`, initialized on lines 289 to 298, are the number of words that exist in both sets, and the number of extra bytes that have to be tested in the potential subset. For example, if set A has 10 bytes and set B has 20, and if you're determining whether set A is a subset of set B, you need to look only at the first 10 bytes of both sets. It doesn't matter whether the last 10 bytes of set B have anything in them. If, however, you want

Deleting sets,
`del_set()`

Duplicating sets,
`dupset()`.

Enlarging sets,
`_addset()`.

Finding the number of
elements in a sets,
`num_ele()`.

Comparing sets for
equivalent, intersecting,
`_set_test()`.

Comparing sets for sort-
ing, `setcmp()`.

Testing for subset,
`subset()`.

Listing A.6. *set.c*— Adding Members to and Enlarging aSET

```

80 PUBLIC int      _addset( set, bit )
81 SET      *set;
82 {
83     /* Addset is called by the ADD() macro when the set isn't big enough. It
84      * expands the set to the necessary size and sets the indicated bit.
85      */
86
87     void enlarge( int, SET* );           /* immediately following */
88
89     enlarge( _ROUND(bit), set );
90     return _GBIT( set, bit, |= );
91 }
92
93 /* -----
94
95 PRIVATE void    enlarge( need, set )
96 SET      *set;
97 {
98     /* Enlarge the set to "need" words, filling in the extra words with zeros.
99      * Print an error message and abort by raising SIGABRT if there's not enough
100     * memory. NULL is returned if raise() returns. Since this routine calls
101     * malloc, it's rather slow and should be avoided if possible.
102     */
103
104     _SETTYPE  *new;
105
106     if( !set || need <= set->nwords )
107         return;
108
109     D( printf("enlarging %d word map to %d words\n", set->nwords, need); )
110
111     if( !(new = (_SETTYPE *) malloc( need * sizeof(_SETTYPE) )) )
112     {
113         fprintf(stderr, "Can't get memory to expand set\n");
114         exit( 1 );
115     }
116     memcpy( new, set->map,           set->nwords * sizeof(_SETTYPE) );
117     memset( new + set->nwords, 0, (need - set->nwords) * sizeof(_SETTYPE) );
118
119     if( set->map != set->defmap )
120         free( set->map );
121
122     set->map    = new ;
123     set->nwords = (unsigned char) need ;
124     set->nbits   = need * _BITS_IN_WORD ;
125 }
```

to know whether the longer set (B) is a subset of the shorter one (A), all bytes that are not in set A (the extra 10 bytes in B) must be scanned to make sure they're all zeros. The common parts of the sets are scanned by the loop on line 303, and the tail of the larger set is scanned, if necessary, by the loop on line 307.

Set operations (union, intersection, symmetric difference, assignment).
`_set_op()`.

The next three functions, in Listing A.8, all modify a set one way or another. The `_set_op()` function on line 313 performs the union, intersection, symmetric difference, and assignment functions.

I've used the same strategy here as I used earlier when the sets were different sizes. The common words are manipulated first, and the tail of the longer set is modified if

Listing A.7. *set.c*— Set-Testing Functions

```

126 PUBLIC int      num_ele( set )
127 SET      *set;
128 {
129     /* Return the number of elements (nonzero bits) in the set. NULL sets are
130     * considered empty. The table-lookup approach used here was suggested to
131     * me by Doug Merrit. Nbits[] is indexed by any number in the range 0-255,
132     * and it evaluates to the number of bits in the number.
133     */
134     static unsigned char nbits[] =
135     {
136         /* 0-15 */ 0, 1, 1, 2, 1, 2, 2, 3, 1, 2, 2, 2, 3, 2, 3, 3, 4,
137         /* 16-31 */ 1, 2, 2, 3, 2, 3, 3, 4, 2, 3, 3, 4, 3, 4, 4, 5,
138         /* 32-47 */ 1, 2, 2, 3, 2, 3, 3, 4, 2, 3, 3, 4, 3, 4, 4, 5,
139         /* 48-63 */ 2, 3, 3, 4, 3, 4, 4, 5, 3, 4, 4, 5, 4, 5, 5, 6,
140         /* 64-79 */ 1, 2, 2, 3, 2, 3, 3, 4, 2, 3, 3, 4, 3, 4, 4, 5,
141         /* 80-95 */ 2, 3, 3, 4, 3, 4, 4, 5, 3, 4, 4, 5, 4, 5, 5, 6,
142         /* 96-111 */ 2, 3, 3, 4, 3, 4, 4, 5, 3, 4, 4, 5, 4, 5, 5, 6,
143         /* 112-127 */ 3, 4, 4, 5, 4, 5, 5, 6, 4, 5, 5, 6, 5, 6, 6, 7,
144         /* 128-143 */ 1, 2, 2, 3, 2, 3, 3, 4, 2, 3, 3, 4, 3, 4, 4, 5,
145         /* 144-159 */ 2, 3, 3, 4, 3, 4, 4, 5, 3, 4, 4, 5, 4, 5, 5, 6,
146         /* 160-175 */ 2, 3, 3, 4, 3, 4, 4, 5, 3, 4, 4, 5, 4, 5, 5, 6,
147         /* 176-191 */ 3, 4, 4, 5, 4, 5, 5, 6, 4, 5, 5, 6, 5, 6, 6, 7,
148         /* 192-207 */ 2, 3, 3, 4, 3, 4, 4, 5, 3, 4, 4, 5, 4, 5, 5, 6,
149         /* 208-223 */ 3, 4, 4, 5, 4, 5, 5, 6, 4, 5, 5, 6, 5, 6, 6, 7,
150         /* 224-239 */ 3, 4, 4, 5, 4, 5, 5, 6, 4, 5, 5, 6, 5, 6, 6, 7,
151         /* 240-255 */ 4, 5, 5, 6, 5, 6, 6, 7, 5, 6, 6, 7, 6, 7, 7, 8
152     };
153     int             i;
154     unsigned int    count = 0;
155     unsigned char   *p;
156
157     if( !set )
158         return 0;
159
160     p = (unsigned char *)set->map ;
161
162     for( i = _BYTES_IN_ARRAY(set->nwords) ; --i >= 0 ; )
163         count += nbits[ *p++ ] ;
164
165     return count;
166 }
167
168 /* ----- */
169
170 PUBLIC int      _set_test( set1, set2 )
171 SET      *set1, *set2;
172 {
173     /* Compares two sets. Returns as follows:
174     *
175     * _SET_EQUIV      Sets are equivalent
176     * _SET_INTER      Sets intersect but aren't equivalent
177     * _SET_DISJ       Sets are disjoint
178     *
179     * The smaller set is made larger if the two sets are different sizes.
180     */
181
182     int             i, rval = _SET_EQUIV ;
183     _SETTYPE      *p1, *p2;
184 }
```



Listing A.7. continued...

```

185     i = max( set1->nwords, set2->nwords );
186
187     enlarge( i, set1 );           /* Make the sets the same size */
188     enlarge( i, set2 );
189
190     p1 = set1->map;
191     p2 = set2->map;
192
193     for( --i >= 0 ; p1++, p2++ )
194     {
195         if( *p1 != *p2 )
196         {
197             /* You get here if the sets aren't equivalent. You can return
198             * immediately if the sets intersect but have to keep going in the
199             * case of disjoint sets (because the sets might actually intersect
200             * at some byte, as yet unseen).
201             */
202
203             if( *p1 & *p2 )
204                 return _SET_INTER ;
205             else
206                 rval = _SET_DISJ ;
207         }
208     }
209
210     return rval;                  /* They're equivalent */
211 }
212
213 /* -----
214
215 PUBLIC setcmp( set1, set2 )
216 SET *set1, *set2;
217 {
218     /* Yet another comparison function. This one works like strcmp(),
219     * returning 0 if the sets are equivalent, <0 if set1<set2 and >0 if
220     * set1>set2.
221     */
222
223     int i, j;
224     _SETTYPE *p1, *p2;
225
226     i = j = min( set1->nwords, set2->nwords );
227
228     for( p1 = set1->map, p2 = set2->map ; --j >= 0 ; p1++, p2++ )
229         if( *p1 != *p2 )
230             return *p1 - *p2;
231
232     /* You get here only if all words that exist in both sets are the same.
233     * Check the tail end of the larger array for all zeros.
234     */
235
236     if( (j = set1->nwords - i) > 0 )                  /* Set 1 is the larger */
237     {
238         while( --j >= 0 )
239             if( *p1++ )
240                 return 1;
241     }

```



Listing A.7. continued...

```
242     else if( (j = set2->nwords - i) > 0)          /* Set 2 is the larger */
243     {
244         while( --j >= 0 )
245             if( *p2++)
246                 return -1;
247     }
248
249     return 0;                                         /* They're equivalent */
250 }
251
252 /* -----
253
254 PUBLIC unsigned sethash( set1 )
255 SET *set1;
256 {
257     /* hash the set by summing together the words in the bit map */
258
259     _SETTYPE      *p;
260     unsigned      total;
261     int          j;
262
263     total = 0;
264     j      = set1->nwords ;
265     p      = set1->map    ;
266
267     while( --j >= 0 )
268         total += *p++;
269
270     return total;
271 }
272
273 /* -----
274
275 PUBLIC int      subset( set, possible_subset )
276 SET      *set, *possible_subset;
277 {
278     /* Return 1 if "possible_subset" is a subset of "set". One is returned if
279     * it's a subset, zero otherwise. Empty sets are subsets of everything.
280     * The routine silently malfunctions if given a NULL set, however. If the
281     * "possible_subset" is larger than the "set", then the extra bytes must
282     * be all zeros.
283     */
284
285     _SETTYPE *subsetp, *setp;
286     int    common;           /* This many bytes in potential subset */
287     int    tail;            /* This many implied 0 bytes in b */
288
289     if( possible_subset->nwords > set->nwords )
290     {
291         common = set->nwords ;
292         tail   = possible_subset->nwords - common ;
293     }
294     else
295     {
296         common = possible_subset->nwords;
297         tail   = 0;
298     }
299 }
```



Listing A.7. continued...

```

300     subsetp = possible_subset->map;
301     setp    = set->map;
302
303     for( --common >= 0; subsetp++, setp++ )
304         if( (*subsetp & *setp) != *subsetp )
305             return 0;
306
307     while( --tail >= 0 )
308         if( *subsetp++ )
309             return 0;
310
311     return 1;
312 }
```

Listing A.8. set.c—Set Manipulation Functions

```

313 PUBLIC void      _set_op( op, dest, src )
314 int          op;
315 SET        *src, *dest;
316 {
317     /* Performs binary operations depending on op:
318     *
319     * _UNION:      dest = union of src and dest
320     * _INTERSECT:   dest = intersection of src and dest
321     * _DIFFERENCE: dest = symmetric difference of src and dest
322     * _ASSIGN:     dest = src;
323     *
324     * The sizes of the destination set is adjusted so that it's the same size
325     * as the source set.
326     */
327
328     _SETTYPE      *d;      /* Pointer to destination map      */
329     _SETTYPE      *s;      /* Pointer to map in set1           */
330     int          ssize;  /* Number of words in src set      */
331     int          tail;   /* dest set is this much bigger */
332
333     ssize = src->nwords ;
334
335     if( (unsigned)dest->nwords < ssize ) /* Make sure dest set is at least */
336         enlarge( ssize, dest );           /* as big as the src set.          */
337
338     tail  = dest->nwords - ssize ;
339     d    = dest->map ;
340     s    = src ->map ;
341
342     switch( op )
343     {
344     case _UNION:      while( --ssize >= 0 )
345             *d++ |= *s++ ;
346             break;
347     case _INTERSECT:  while( --ssize >= 0 )
348             *d++ &= *s++ ;
349             while( --tail >= 0 )
350                 *d++ = 0;
351             break;
352 }
```



Listing A.8. continued...

```

352     case _DIFFERENCE: while( --ssize >= 0 )
353         *d++ ^= *s++ ;
354         break;
355     case _ASSIGN:    while( --ssize >= 0 )
356         *d++ = *s++ ;
357         while( --tail >= 0 )
358             *d++ = 0;
359         break;
360     }
361 }
362 */
363 /* -----
364 PUBLIC void      invert( set )
365 SET      *set;
366 {
367     /* Physically invert the bits in the set. Compare with the COMPLEMENT()
368     * macro, which just modifies the complement bit.
369     */
370
371     _SETTYPE *p, *end ;
372
373     for( p = set->map, end = p + set->nwords ; p < end ; p++ )
374         *p = ~*p;
375     }
376 */
377 /* -----
378 PUBLIC void      truncate( set )
379 SET      *set;
380 {
381     /* Clears the set but also sets it back to the original, default size.
382     * Compare this routine to the CLEAR() macro which clears all the bits in
383     * the map but doesn't modify the size.
384     */
385
386     if( set->map != set->defmap )
387     {
388         free( set->map );
389         set->map = set->defmap;
390     }
391     set->nwords = _DEFWORDS;
392     set->nbits  = _DEFBITS;
393     memset( set->defmap, 0, sizeof(set->defmap) );
394 }
395 */
396 }
```

necessary. The work is all done by the **while** loops in the **switch** on lines 342 to 360. It's probably better style to put one **while** statement outside the switch than to put several identical ones in the **cases**, but the latter is more efficient because the **switch** won't have to be re-evaluated on every iteration of the loop. The first **while** loop in every case takes care of all destination elements that correspond to source elements. The maps are processed one word at a time, ORing the words together for union, ANDing them for intersection, XORing them for symmetric difference, and just copying them for assignment. Note that the source set can not be larger than the destination set because of the earlier `enlarge()` call on line 336.

The second **while** loop in each **case** won't execute if the two sets are the same size, because the tail size will be zero. Otherwise, the destination set is larger and you must do different things to the tail, depending on the operation. Since the missing elements of the smaller set are all implied zeros, the following operations are performed:

- union: Do nothing else to the destination because there are no more bits in the source set to add.
- intersection: Clear all the bits in the tail of the destination because no bits are set in the source so there's no possibility of intersection.
- difference: Do nothing because every bit that's set in the source set will not be set in the destination.
- assignment: Set all bits in the tail of the destination to 0 (because the implied bits in the source are all 0).

The `invert()` subroutine starting on line 365 of Listing A.8 just goes through the map, reversing the sense of the bits. The `truncate()` function on line 380 restores a set to its initial, empty, condition. This last routine is really a more efficient replacement for:

```
delset( s );
s = newset();
```

You may be better off—in terms of speed—to clear the existing set with `CLEAR` rather than calling `truncate()`, because `free()` is pretty slow.

Accessing an entire set,
`next_member()`.

The final two set routines, which access and print entire sets, are in Listing A.9. The `next_member()` function on line 397 lets you access all elements of a set sequentially. When the function is called several successive times with the same argument, it returns the next element of the set with each call, or `-1` if there are no more elements. Every time the `set` argument changes, the search for elements starts back at the beginning of the set. Similarly, `next_member(NULL)` resets the search to the beginning of the set (and does nothing else). You should not put any new elements in the set between a `next_member(NULL)` call and a subsequent `next_member(set)` call. Elements should not be added to the set between successive `next_member()` calls.

Listing A.9. *set.c*—Getting Elements and Printing the Set

```
397 PUBLIC int      next_member( set )
398 SET      *set;
399 {
400     /* set == NULL                      Reset
401      * set changed from last call:    Reset and return first element
402      * otherwise                     return next element or -1 if none.
403     */
404
405     static SET  *oset      = NULL;   /* "set" arg in last call      */
406     static int   current_member = 0;   /* last-accessed member of cur. set */
407     _SETTYPE   *map;
408
409     if( !set )
410         return( (int)( oset = NULL ) );
411
412     if( oset != set )
413     {
414         oset      = set;
415         current_member = 0 ;
416     }
```



Listing A.9. continued...

```

417     for(map = set->map; *map == 0 && current_member < set->nbits; ++map)
418         current_member += _BITS_IN_WORD;
419     }
420
421     /* The increment must be put into the test because, if the TEST() invocation
422      * evaluates true, then an increment on the right of a for() statement
423      * would never be executed.
424      */
425
426     while( current_member++ < set->nbits )
427         if( TEST(set, current_member-1) )
428             return( current_member-1 );
429     return( -1 );
430 }
431
432 /* -----
433
434 PUBLIC void      pset( set, output_routine, param )
435 SET      *set;
436 int       (*output_routine)();
437 void      *param;
438 {
439     /* Print the contents of the set bit map in human-readable form. The
440      * output routine is called for each element of the set with the following
441      * arguments:
442      *
443      * (*out)( param, "null",      -1);  Null set ("set" arg == NULL)
444      * (*out)( param, "empty",     -2);  Empty set (no elements)
445      * (*out)( param, "%d ",      N);   N is an element of the set
446      */
447
448     int i, did_something = 0;
449
450     if( !set )
451         (*output_routine)( param, "null", -1 );
452     else
453     {
454         next_member( NULL );
455         while( (i = next_member(set))  >= 0 )
456         {
457             did_something++;
458             (*output_routine )( param, "%d ", i );
459         }
460         next_member( NULL );
461
462         if( !did_something )
463             (*output_routine )(param, "empty", -2 );
464     }
465 }
```

A major limitation to `next_member()` is that calls to `next_member` on different sets cannot be interspersed. This problem could be solved by putting the `current_member` counter (declared on line 406) into the `SET` structure rather than keeping it as a local `static` variable. Since I've never needed this functionality, I didn't want to add the extra overhead needed to support it. The loop on line 417 is skipping initial words in the bit map that are all zeros.

Printing a set, pset().

The pset() function on line 434 prints all the elements in a set. The standard call:

```
pset( set, fprintf, stdout );
```

prints the elements of the set, separated by space characters, and without a newline at the end of the list. The second argument can actually be a pointer to any function, however. The third argument is just passed to that function, along with a format string. The function is called indirectly through a pointer as follows:

```
(*out)( param, "null", -1);    For null sets (set is NULL)
(*out)( param, "empty", -2);   For empty set (set has no elements)
(*out)( param, "%d ", N);     Normally, N is an element of the set
```

Calls to pset() and next_member() should not be interspersed.

A.3 Database Maintenance—Hashing:

Compilers all need some sort of simple data-base management system. Typically the databases are small enough to fit into memory, so an elaborate file-based system is not only not required, but is contraindicated (because of the excessive access time). Of the various techniques that are available for data management, the most appropriate is the hash table.

The eight functions in this section implement a general-purpose, data-base manager that uses a hash strategy, but you can use these functions without any knowledge of the mechanics of manipulating a hash table. The hash functions are used in Chapter Two, but hashing isn't discussed until Chapter Six. You may want to read that discussion before proceeding. You can also just skim over the the function overviews that follow, skipping over the implementation details until you've read Chapter Six.

Listing A.10 shows a very simple application that creates a database that holds the argv entries, and then prints the entries. The details of all the function calls will be discussed in a moment, but to summarize: A single database record is defined with the **typedef** on line nine. An empty database is created on line 55 with a maketab() call, which is passed pointers to two auxiliary functions that control database manipulation: The hash function on line 17 translates the key field to a pseudo-random number, here using a shift and exclusive-OR strategy. The comparison function on line 31 does a lexicographic comparison of the key fields of two database records. The newsym() call on line 58 allocates space for a new database record, which is put into the database with the addsym() call on line 60.

Listing A.10. Program to Demonstrate Hash Functions

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/hash.h>
4
5 /* A application that demonstrates how to use the basic hash functions. Creates
6 * a database holding the argv strings and then prints the database.
7 */
8
9 typedef struct                                /* A database record */
10 {
11     char *key;
12     int other_stuff;
13 } ENTRY;
14

```



Listing A.10. continued...

```

15  /*-----*/
16
17 unsigned hash( sym )                                /* Hash function. Convert the key */
18 ENTRY    *sym;                                     /* to a number. */
19 {
20     char *p;
21     int hash_value = 0;
22
23     for( p = sym->key; *p ; hash_value = (hash_value << 1) ^ *p++ )
24     ;
25
26     return hash_value;
27 }
28
29 /*-----*/
30
31 int cmp( sym1, sym2 )                            /* Compare two database records. */
32 ENTRY    *sym1, *sym2;                           /* Works like strcmp(). */
33 {
34     return strcmp( sym1->key, sym2->key );
35 }
36
37 /*-----*/
38
39 void print( sym, stream )                      /* print a database record to the stream */
40 ENTRY    *sym;
41 FILE    *stream;
42 {
43     fprintf( stream, "%s\n", sym->key );
44 }
45
46 /*-----*/
47
48 main( argc, argv )
49 int      argc;
50 char    **argv;
51 {
52     HASH_TAB *tab;
53     ENTRY    *p;
54
55     tab = maketab(31, hash, cmp);                /* make hash table */
56     for( ++argv, --argc; --argc>=0; ++argv)    /* For each element of argv */
57     {
58         p = (ENTRY *) newsym(sizeof(ENTRY));    /* put it into the table */
59         p->key = *argv;
60         addsym( tab, p );
61     }
62                                         /* print the table. stdout is */
63     ptab( tab, print, stdout, 1 );               /* passed through to print(). */
64 }
```

Create hash table.

```
#include <tools/hash.h>
HASH_TAB *maketab      (unsigned maxsym, unsigned (*hash) (),
                           unsigned (*cmp) ())
```

Make a hash table of the size specified in maxsym. The hash table is a data structure that the manager uses to organize the database entries. It contains a “hash table” array along with various other housekeeping variables. The maxsym argument controls the size of the array component of the data structure, but any number of entries can be put into the table, regardless of the table size. Different table sizes will affect the speed with which a record can be accessed, however. If the table is too small, the search time will be unnecessarily long. Ideally the table should be about the same size as the expected number of table entries. There is no benefit in making the table too large.³ It’s a good idea to make maxsym a prime number⁴ Some useful sizes are: 47, 61, 89, 113, 127, 157, 193, 211, 257, 293, 337, 367, 401. If maxsym is zero, 127 is used.

The functions referenced by the two pointer arguments (hash and cmp) are used to manipulate the database. The hash function is called indirectly, like this:

```
(*hash) ( sym );
```

where sym is a pointer to a region of memory allocated by newsym(), which is used like malloc() to allocate a new table element. The assumption is that newsym() is getting space for a structure, one field of which is the key. The hash function should return a pseudo-random number, the value of which is controlled by the key; the same key should always generate the same number, but different keys should generate different numbers. The simplest, but by no means the best, hash strategy just adds together the characters in the name. Better methods are discussed in Chapter Six. Two default hash functions are discussed below.

The comparison function (cmp) is passed two pointers to database records. It should compare the key fields and return a value representing the ordering of the keys in a manner analogous to strcmp(). A call to (*cmp)(p1, p2) should return a negative number if the key field in the structure pointed to by p1 is less than the one in *p2. It should return 0 if the two keys are identical, and it should return a positive number if the key in *p2 is greater.

maketab() prints an error message and raises SIGABRT if there’s not enough memory. (It works the same way as newset() in this regard.)

Allocate memory for symbol.

```
void *newsym(int size)
```

This routine is used like malloc() to get space for a database record. The returned memory is initialized to zeros. Typically, you use newsym() to allocate a structure, one field of which is the key. The routine prints an error message and raises SIGABRT if there’s not enough memory. The pointer returned from newsym() may not be passed to free(); use freesym(), below.

Free memory used by symbol.

```
void freesym(void *sym)
```

This routine frees the memory for a symbol created by a previous newsym()

3. You tend to get no fewer collisions in a too-long table. You just get holes.

4. The distribution of elements in the table tends to be better if the table size is prime.

call. You may not use `free()` for this purpose. Do not free symbols that are still in the table—remove them with a `delsym()` call first.

void *`addsym(HASH_TAB *tabp, void *sym)`

Add symbol to table.

Add a symbol to the hash table pointed to by `tabp`—a pointer returned from a previous `maketab()` call. The `sym` argument points at a database record, a pointer to which was returned from a previous `newsym()` call. You must initialize the key field of that record prior to the `addsym()` call.

void *`findsym(HASH_TAB *tabp, void *sym)`

Find symbol in table.

Return either a pointer to a previously-inserted database record or `NULL` if the record isn't in the database. If more than one entry for a given key is in the database, the most recently added one is found. The `sym` argument is used to identify the record for which you're searching. It is not used directly by `findsym()`, but is passed to the hash and comparison functions. The comparison function is called as follows:

```
(*cmp)(sym, item_in_table);
```

Here, `item_in_table` is a pointer to an arbitrary database element, and `sym` is just the second argument to `findsym()`.

Strictly speaking, `sym` should be a pointer to an initialized database record returned from `newsym()`. It's inconvenient, though, to allocate and initialize a structure just to pass it to `findsym()`. You can get around the problem in one common situation. If the key field is a character array, and that array is the first field in the structure, you can pass a character-string name to `findsym()` as the key. This is a hack, but is nonetheless useful. The technique is illustrated in Listing A.11. This technique works only if the key is a character array—the string must be physically present as the first few bytes of the structure. Character pointers won't work. Note that `strcmp()` is used as the comparison function. This works only because the array is at the top of the structure, so the structure pointer passed to `strcmp` is also the address of the key array.

Simplifying `findsym()` calls by subterfuge.

void `delsym(HASH_TAB *tabp, void *sym)`

Remove symbol from table.

Remove a symbol from the hash table, a pointer to which was returned by a previous `findsym()` call. `sym` is the pointer returned from `findsym()` and `tabp` is a pointer returned from `maketab()`. The record is removed from the table, but the associated memory is not freed, so you can recycle it—you can reinitialize the record and reinsert it into the table at a later time. Use `freesym()` to free the memory used by `sym`—you may not use `free()` for this purpose. It's a serious error to delete a symbol that isn't actually in the table—`addsym()` must have been called for a node before `delsym()` can be called.

void *`nextsym(HASH_TAB *tabp, void *last)`

Get next symbol.

This function finds all references to objects in a table that have the same name. The first such object is found by `find_sym()`. The second object is found by passing the pointer returned from `find_sym()` to `nextsym()`, which returns either a pointer to the next object or `NULL` if there are no such objects. Use it like this:

Listing A.11. Fooling findsym()

```

1  typedef struct
2  {
3      char    name[ SIZE ];           /* must be first, and must be an array */
4      int     other_stuff;
5  } ENTRY;
6
7  hash( key )
8  char *key;
9  {
10     int   i = 0;                /* Add together characters in the name      */
11     while( *key )             /* use a left shift to randomize the number */
12         i += *key++ << 1 ;
13 }
14
15 ignaz()
16 {
17     HASH_TAB    *tab;
18     extern     hash_add();      /* hash function */
19     extern     strcmp();
20     ENTRY       *p;
21     . . .
22     tab = maketab( 61, hash, strcmp );
23     . . .
24     p = findsym( Tab, "krazy" );
25 }
```

```

        p = findsym( Tab, "krazy" );    /* Get the first one. */
        . . .
if( !(p = nextsym(Tab, p)) )    /* Get the next one. */
        /* no more symbols */
```

Third and subsequent objects are found by passing `nextsym()` the value returned from the previous `nextsym()` call.

Print entire database.

```
int ptab(HASH_TAB *tabp, void (*print)(), void *param,
          int sort)
```

Print all records in the database represented by the hash table pointed to by `tabp`. The function pointed to by `print` is called for every element of the table as follows:

```
(*print)( sym, param )
```

If `sort` is false, the table elements are printed in random order and 1 is always returned. If `sort` is true, the table is printed only if the routine can get memory to sort the table. Zero is returned (and nothing is printed) if memory isn't available, otherwise 1 is returned and the table is printed in order controlled by the comparison function passed to `maketab()`. If this comparison function works as described earlier, the table is printed in ascending order. Reverse the sense of the return value to print in descending order. In the current example, you can change the comparison function as follows, reversing the arguments to `strcmp()`:

```

    cmp( a, b )                      /* Print in descending order */
    ENTRY *a, *b;
{
    return strcmp( b->key, a->key );
}

unsigned hash_add(char *name);
unsigned hash_pjw(char *name);

```

Default hash functions.

These two functions are hash functions that you can pass to `maketab()`. They are passed character strings and return a pseudo-random integer generated from that string. `hash_add()` just adds the character in the name—it's fast but doesn't work well if the table size is larger than 128 or if keys are likely to be permutations of each other. The `hash_pjw()` function uses a shift-and-exclusive-OR algorithm that yields better results at the cost of execution speed. As with `findsym()`, if the table entries have a character-array key at the top of the structure, these functions can be used directly by `maketab()`,

```

typedef struct
{
    char key[ 80 ];
    int stuff;
}
ENTRY;
...
maketab(..., hash_pjw, ...);

```

Otherwise, you'll have to encapsulate the hash function inside a second function like this:

```

typedef struct
{
    int stuff;
    char *key;
}
ENTRY;

hash_funct( sym )
ENTRY      *sym;
{
    return hash_add( sym->key );
}
...
maketab(..., hash_funct, ...);

```

A.3.1 Hashing—Implementation

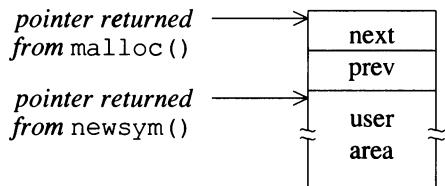
I've used data abstraction extensively in the package presented here. That is, the mechanics of manipulating the database is hidden from the user of the routines. For example, the user calls one function to allocate a data-base block of an arbitrary size in a manner similar to `malloc()`. The block is then modified as necessary, and it's inserted into the database with a second function call. The mechanics of allocating space, of doing the insertion, and so forth, are hidden in the functions. Similarly, the internal data structures used for database maintenance are also hidden. This abstraction makes these routines quite flexible. Not only can you use them in disparate applications, but you can change the way that the database is maintained. As long as the subroutine-call interface is the same, you don't have to change anything in the application program.

Data abstraction.

The hash functions work with two data structures. I'll look at the C definitions and allocation functions first, and then describe how they work.

A hash-table element,
BUCKET, newsym().

The first structure is the BUCKET, declared in *hash.h*, (the start of which is in Listing A.12). The newsym() function, (in Listing A.12) allocates a BUCKET along with the memory requested by the user. newsym() simultaneously allocates enough memory for both the BUCKET header, and for a user space, the size of which is passed in as a parameter. It returns a pointer to the area just below the header, which can be used in any way by the application program:



Listing A.12. *hash.h*— The BUCKET

```

1  typedef struct BUCKET
2  {
3      struct BUCKET      *next;
4      struct BUCKET      **prev;
5
6  } BUCKET;
  
```

Listing A.13. *hash.c*— BUCKET Allocation

```

9  PUBLIC void *newsym( size )
10 int      size;
11 {
12     /* Allocate space for a new symbol; return a pointer to the user space. */
13
14     BUCKET      *sym;
15
16     if( !(sym = (BUCKET *) calloc( size + sizeof(BUCKET), 1)) )
17     {
18         fprintf( stderr, "Can't get memory for BUCKET\n" );
19         raise( SIGABRT );
20         return NULL;
21     }
22     return (void *)( sym + 1 );           /* return pointer to user space */
23 }
24
25 /-----*/
26
27 PUBLIC void freesym( sym )
28 void      *sym;
29 {
30     free( (BUCKET *)sym - 1 );
31 }
  
```

Freeing BUCKETS,
freesym().

The freesym() function, also in Listing A.13, frees the memory allocated by a previous newsym(). It just backs up the sym pointer to its original position and then calls free().

The hash table itself,
HASH_TAB.

The second data structure of interest is the HASH_TAB which holds the actual hash table itself. (It's declared in Listing A.14.) Like a BUCKET, it's a structure of variable length. The header contains the table size in elements (size), the number of entries

currently in the table (numsyms), pointers to the hash and comparison functions (hash and cmp), and the table itself (table is the first element of the table). The numsyms field is used only for statistical purposes—it's not needed by any of the hash-table functions.

The maketab() function (in Listing A.15) allocates a single chunk of memory big enough to hold both the header and an additional area that will be used as the array. The table is declared as a one-element array, but the array can actually be any size, provided that there's enough available memory following the header. I'm taking advantage of the fact that C doesn't do array-boundary checking when the array is accessed.

Making a hash table,
maketab().

Listing A.14. *hash.h*—HASH_TAB Definition

```

7  typedef struct  hash_tab_
8  {
9      int          size      ;           /* Max number of elements in table      */
10     int          numsyms   ;           /* number of elements currently in table */
11     unsigned    (*hash) () ;           /* hash function                         */
12     int          (*cmp) () ;           /* comparison funct, cmp(name,bucket_p); */
13     BUCKET      *table[1];           /* First element of actual hash table    */
14
15 } HASH_TAB;

```

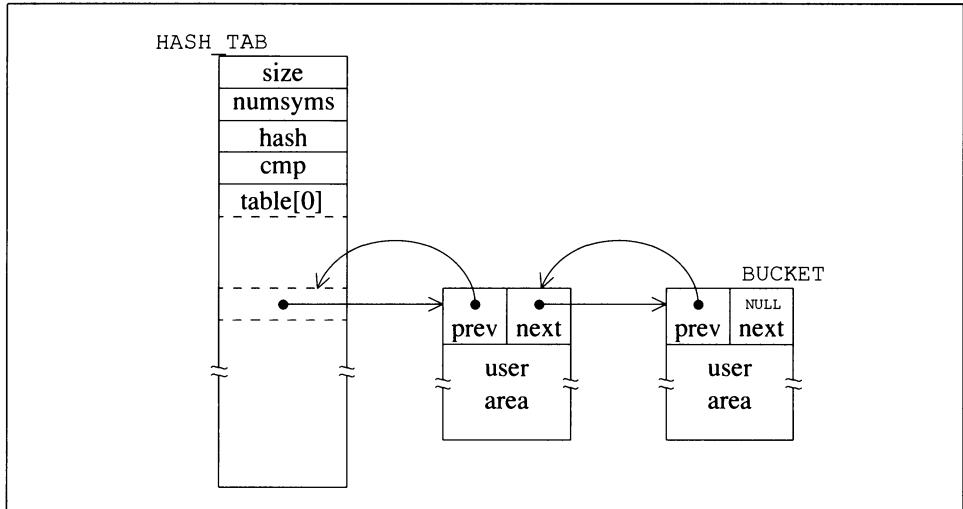
Listing A.15. *hash.c*—HASH_TAB Allocation

```

32 PUBLIC HASH_TAB *maketab( maxsym, hash_function, cmp_function )
33 unsigned          maxsym;
34 unsigned          (*hash_function)();
35 int              (*cmp_function)();
36 {
37     /* Make a hash table of the indicated size. */
38
39     HASH_TAB *p;
40
41     if( !maxsym )
42         maxsym = 127;
43         /* |<-- space for table ---->|<- and header -->| */
44     if( p=(HASH_TAB*) calloc(1, (maxsym * sizeof(BUCKET*)) + sizeof(HASH_TAB)) )
45     {
46         p->size      = maxsym      ;
47         p->numsyms   = 0          ;
48         p->hash      = hash_function ;
49         p->cmp       = cmp_function ;
50     }
51     else
52     {
53         fprintf(stderr, "Insufficient memory for symbol table\n");
54         raise( SIGABRT );
55         return NULL;
56     }
57     return p;
58 }

```

These two structures interact as shown in Figure A.1. The actual table is an array of BUCKET pointers. Empty table elements are set to NULL, and new BUCKETS are tacked onto the head of the list. BUCKETS in a HASH_TAB.

Figure A.1. A Hash Table with Two Elements in It

Adding symbols to the table, `addsym()`.

The `addsym()` function, which puts items into the table, is in Listing A.16. A pointer to the current hash-table element, `p`, is initialized on line 68 by calling the hash function indirectly through the `hash` pointer in the `HASH_TAB` header. Note that the `BUCKET` pointer (`sym`) comes into the routine pointing at the `user area`. It is decremented to point at the `BUCKET` header on line 68 after the hash function is called. The code on lines 70 to 73 link the new node to the head of a chain found at the previously computed array element.

Listing A.16. *hash.c*— Adding a Symbol to the Table

```

59 PUBLIC void *addsym( tabp, isym )
60 HASH_TAB *tabp;
61 void      *isym;
62 {
63     /* Add a symbol to the hash table. */
64
65     BUCKET    **p, *tmp ;
66     BUCKET    *sym = (BUCKET *)isym;
67
68     p = & (tabp->table)[ (*tabp->hash)( sym-- ) % tabp->size ];
69
70     tmp      = *p  ;
71     *p      = sym ;
72     sym->prev = p  ;
73     sym->next = tmp ;
74
75     if( tmp )
76         tmp->prev = &sym->next ;
77
78     tabp->numsyms++;
79     return (void *) (sym + 1);
80 }
```

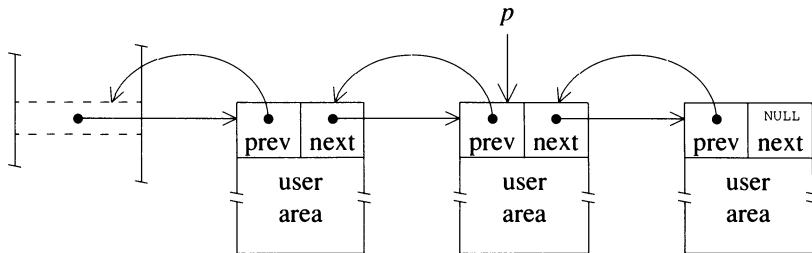
BUCKETS form a doubly linked list, backwards pointers.

Note that the chain of BUCKETS is a doubly-linked list. You need the backwards pointers to delete an arbitrary element in the table without having to search for that element. The only obscure point is the two stars in the definition of the backwards

pointer—the forward pointer (`next`) is a pointer to a BUCKET, but the backwards pointer (`prev`) is a pointer to a BUCKET pointer. You need this extra level of indirection because the head of the chain is a simple pointer, not an entire BUCKET structure. The backwards pointer for the leftmost node in the chain points at the head-of-chain pointer. All other backwards pointers hold the address of the `next` field from the previous node.

You can see the utility of this system by looking at the code necessary to delete an arbitrary node. Say that you want to delete the node pointed to by `p` in the following picture:

Deleting a BUCKET from the table, `delsym()`.



The removal can be accomplished with the following statement:

```
if( *(p->prev) = p->next )
    p->next->prev = p->prev ;
```

The pointer from the previous node (the one that points at the node to delete) is modified first so that it points around the deleted node. Then the backwards pointer from the next node is adjusted. The `if` is required because the `next` pointer is `NULL` on the last node in the chain. The double indirection on the backwards pointer makes this code work regardless of the position of the node in the chain—the first and last nodes are not special cases. The `delsym()` function, which removes an arbitrary node from the table, is shown in Listing A.17.

Listing A.17. *hash.c*—Removing A Node From The Table

```
81 PUBLIC void delsym( tabp, isym )
82 HASH_TAB *tabp;
83 void *isym;
84 {
85     /* Remove a symbol from the hash table. "sym" is a pointer returned from
86      * a previous findsym() call. It points initially at the user space, but
87      * is decremented to get at the BUCKET header.
88 */
89
90     BUCKET *sym = (BUCKET *)isym;
91
92     if( tabp && sym )
93     {
94         --tabp->numpms;
95         --sym;
96
97         if( *(sym->prev) = sym->next )
98             sym->next->prev = sym->prev ;
99     }
100 }
```

Finding symbols,
findsym(), nextsym().

The two symbol-finding functions, findsym() and nextsym(), are in Listing A.18. findsym() just hashes to the correct place in the table and then chases down the chain, looking for the required node. It returns a pointer to the user area of the BUCKET (thus the +1 on line 119), or NULL if it can't find the required node. nextsym() just continues chasing down the same chain, starting where the last search left off.

Listing A.18. hash.c—Finding a Symbol

```

101 PUBLIC void      *findsym( tabp, sym )
102 HASH_TAB        *tabp;
103 void            *sym;
104 {
105     /*  Return a pointer to the hash table element having a particular name
106     *  or NULL if the name isn't in the table.
107     */
108
109     BUCKET        *p ;
110
111     if( !tabp )           /* Table empty */
112         return NULL;
113
114     p = (tabp->table)[ (*tabp->hash)(sym) % tabp->size ];
115
116     while( p  &&  (*tabp->cmp)( sym, p+1 )  )
117         p = p->next;
118
119     return (void *) ( p ? p + 1 : NULL );
120 }
121
122 /*-----*/
123
124 PUBLIC void      *nextsym( tabp, i_last )
125 HASH_TAB        *tabp;
126 void            *i_last;
127 {
128     /* Return a pointer the next node in the current chain that has the same
129     * key as the last node found (or NULL if there is no such node). "last"
130     * is a pointer returned from a previous findsym() or nextsym() call.
131     */
132
133     BUCKET *last = (BUCKET *) i_last;
134
135     for(--last; last->next ; last = last->next )
136         if( (tabp->cmp)(last+1, last->next +1) == 0 )    /* keys match */
137             return (char *) ( last->next + 1 );
138
139     return NULL;
}

```

Printing the entire data-
base:ptab().

The last of the support functions is ptab(), which prints the table. It starts on line 142 of Listing A.19. The loop on lines 166 to 173 prints the table in the most straightforward manner. The outer loop goes through the table from top to bottom looking for chains. The inner loop traverses the chain, calling the print function indirectly at each node.

The **else** clause on lines 177 to 214 handles sorted arrays. It allocates an array of BUCKET pointers on line 184 and initializes the array to point at every BUCKET in the table with the loop starting on line 189. (This is essentially the same loop that printed the unsorted array). The new array is sorted on line 208 using assort(), a variant of

Listing A.19. *hash.c*—Printing the Table

```

140 PRIVATE int      (*User_cmp) ();
141
142 PUBLIC int ptab( tabp, print, param, sort )
143 HASH_TAB          *tabp;           /* Pointer to the table */
144 void              (* print)();    /* Print function used for output */
145 void              *param;         /* Parameter passed to print function */
146 int               sort;          /* Sort the table if true. */
147 {
148     /* Return 0 if a sorted table can't be printed because of insufficient
149     * memory, else return 1 if the table was printed. The print function
150     * is called with two arguments:
151     *     (*print)( sym, param )
152     *
153     * Sym is a pointer to a BUCKET user area and param is the third
154     * argument to ptab.
155     */
156
157     BUCKET      **outtab, **outp, *sym, **syntab ;
158     int          internal_cmp();
159     int          i;
160
161     if( !tabp || tabp->size == 0 )    /* Table is empty */
162     return 1;
163
164     if( !sort )
165     {
166         for( syntab = tabp->table, i = tabp->size ; --i >= 0 ; syntab++ )
167         {
168             /* Print all symbols in the current chain. The +1 in the print call
169             * increments the pointer to the applications area of the bucket.
170             */
171             for( sym = *syntab ; sym ; sym = sym->next )
172                 (*print)( sym+1, param );
173         }
174     }
175     else
176     {
177         /* Allocate memory for the outtab, an array of pointers to
178         * BUCKETS, and initialize it. The outtab is different from
179         * the actual hash table in that every outtab element points
180         * to a single BUCKET structure, rather than to a linked list
181         * of them.
182         */
183
184         if( !( outtab = (BUCKET **) malloc(tabp->numsyms * sizeof(BUCKET*)) ) )
185         return 0;
186
187         outp = outtab;
188
189         for( syntab = tabp->table, i = tabp->size ; --i >= 0 ; syntab++ )
190             for( sym = *syntab ; sym ; sym = sym->next )
191             {
192                 if( outp > outtab + tabp->numsyms )
193                 {
194                     fprintf(stderr,"Internal error [ptab], table overflow\n");
195                     exit(1);
196                 }
197             }

```



Listing A.19. continued...

```

198             *outp++ = sym;
199         }
200
201     /*      Sort the outtab and then print it. The (*outp)+1 in the
202     *      print call increments the pointer past the header part
203     *      of the BUCKET structure. During sorting, the increment
204     *      is done in internal_cmp.
205     */
206
207     User_cmp = tabp->cmp;
208     assort( outtab, tabp->numpssyms, sizeof( BUCKET* ), internal_cmp );
209
210     for( outp = outtab, i = tabp->numpssyms; --i >= 0 ; outp++ )
211         (*print)( (*outp)+1, param );
212
213     free( outtab );
214 }
215 return 1;
216 }
217
218 PRIVATE int internal_cmp( p1, p2 )
219 BUCKET **p1, **p2;
220 {
221     (*User_cmp)( *p1 + 1, *p2 + 1 );
222 }
```

the UNIX `qsort()` function that uses a Shell sort rather than a quicksort. (It's discussed later on in this appendix.) The sorted array is printed by the loop on line 210.

The sorting is complicated because, though the comparison function must be passed two pointers to the user areas of buckets, `assort()` passes the sort function two pointers to array elements. That is, `assort()` passes two pointers to BUCKET pointers to the comparison function. The problem is solved by putting a layer around the user-supplied comparison function. `assort()` calls `internal_cmp()`, declared on line 218, which strips off one level of indirection and adjusts the BUCKET pointer to point at the user space before calling the user-supplied function. The user's comparison functions must be passed by a global variable (`User_cmp`, declared on line 140).

Finishing up with the hash-table functions themselves, remainder of `hash.h` contains various function prototypes—it's shown in Listing A.20. Similarly, there are a bunch of `#includes` at the top of `hash.c`, shown in Listing A.21.

Listing A.20. hash.h—HASH_TAB Function Prototypes

```

16 extern HASH_TAB *maketab P(( unsigned maxsym, unsigned (*hash)(), int(*cmp)()));
17 extern void      *newsym  P(( int size ));
18 extern void      freesym  P(( void *sym ));
19 extern void      *addsym  P(( HASH_TAB *tabp, void *sym ));
20 extern void      *findsym  P(( HASH_TAB *tabp, void *sym ));
21 extern void      *nextsym  P(( HASH_TAB *tabp, void *last ));
22 extern void      delsym   P(( HASH_TAB *tabp, void *sym ));
23 extern int       ptab     P(( HASH_TAB *tabp, void(*prnt)(), void *par, int srt));
24 unsigned hash_add    P(( unsigned char *name ));           /* in hashadd.c */
25 unsigned hash_pjw    P(( unsigned char *name ));           /* in hashpjw.c */
```

Listing A.21. *hash.c*—**#includes**

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <signal.h>
4 #include <stdlib.h>
5 #include <string.h>
6 #include <tools/debug.h>
7 #include <tools/hash.h>
```

A.3.2 Two Hash Functions

Implementations of the two default hash functions discussed in Chapter Six are shown in Listings A.22 and A.23.

Listing A.22. *hashadd.c*—An Addition-Based Hash Function

```

1 #include <tools/debug.h>
2 #include <tools/hash.h>
3 /*-
4  * Hash function for use with the functions in hash.c. Just adds together
5  * characters in the name.
6  */
7
8 unsigned hash_add( name )
9 unsigned char *name;
10 {
11     unsigned h ;
12
13     for( h = 0; *name ; h += *name++ )
14         ;
15     return h;
16 }
```

I've modified Aho's version of *hash_pjw* considerably in order to make it portable—the original version assumed that the target machine had a 32-bit **unsigned int**. All the macros on lines four to seven of Listing A.23 are for this purpose. **NBITS_IN_UNSIGNED** evaluates to the number of bits in an **unsigned int**, using the **NBITS** macro from *<tools/debug.h>*. You could also use the ANSI **CHAR_BIT** macro (defined in *limits.h*) and multiply it by eight, but that's risky because a **char** might not be eight bits wide. **SEVENTY_FIVE_PERCENT** evaluates to the number of bits required to isolate the bottom $\frac{3}{4}$ of the number. Given a 16-bit **int**, it will evaluate to 12. **TWELVE_PERCENT** works the same way, but it gets $\frac{1}{8}$ of the number—2, given a 16-bit **int**. **HIGH_BITS** is a mask that isolates the bits in the top $\frac{1}{8}$ of the number—it's 0x3ff for a 16-bit **int**.

Implementing
hash_pjw().

Note that most of the computation in the macros is folded out of existence by the optimizer (because everything's a constant, so the arithmetic can be performed at compile time). For example:

```
h ^= g >> (int)((sizeof(unsigned) * CHAR_BIT) * .75 );
```

generates the following code with Microsoft C, version 5.0:

Listing A.23. *hashpjw.c*—The *hashpjw* Function

```

1 #include <tools/debug.h>
2 #include <tools/hash.h>           /* for prototypes only */
3
4 #define NBITS_IN_UNSIGNED      (      NBITS(unsigned int)      )
5 #define SEVENTY_FIVE_PERCENT   ((int)( NBITS_IN_UNSIGNED * .75 ))
6 #define TWELVE_PERCENT        ((int)( NBITS_IN_UNSIGNED * .125 ))
7 #define HIGH_BITS              ( ~( (unsigned)(~0) >> TWELVE_PERCENT) )
8
9 unsigned hash_pjw( name )
10 unsigned char *name;
11 {
12     unsigned h = 0;                  /* Hash value */
13     unsigned g;
14
15     for( ; *name ; ++name )
16     {
17         h = (h << TWELVE_PERCENT) + *name ;
18         if( g = h & HIGH_BITS )
19             h = (h ^ (g >> SEVENTY_FIVE_PERCENT)) & ~HIGH_BITS ;
20     }
21     return h;
22 }
```

```

    mov     cl,12
    shr     ax,cl
    xor     dx,ax
```

Similarly:

```
g = h & ~( (unsigned)(~0) >> ((int)(NBITS_IN_UNSIGNED * .125)))
```

becomes:

```

    mov     dx,ax
    and     ax,-16384
```

If you modify these constants, don't convert the floating-point numbers to fractions, or else they'll end up as zero. For example, (NBITS_IN_UNSIGNED * (1/8)) won't work because the 1/8 will evaluate to 0 (integer arithmetic is used).

The algorithm uses a shift-and-XOR strategy to randomize the input key. The main iteration of the loop shifts the accumulated hash value to the left by a few bits and adds in the current character. When the number gets too large, it is randomized by XORing it with a shifted version of itself.

A.4 The ANSI Variable-Argument Mechanism

This book uses the ANSI-approved method of supporting subroutines with a variable number of arguments, which uses various macros declared in *<stdarg.h>*, for this purpose. Since not all C environments support *<stdarg.h>*, this section presents a version of the macros in this file. Before looking at *<stdarg.h>* itself, however, the following subroutine demonstrates how the macros are used. The *print_int* subroutine is passed an argument count followed by an arbitrary number of *int*-sized arguments, and it prints those arguments.⁵

Using the *<stdarg.h>*
macros.

```

print_int( arg_count, ... )
int arg_count;
{
    va_list args;

    va_start( args, arg_count );

    while( --arg_count >= 0 )
        printf( "%d ", va_arg(args, int) );

    va_end( args );
}

```

The `args` variable is a pointer to the argument list. It is initialized to point at the second argument in the list with the `va_start()` call, which is passed the argument pointer (`args`) and the name of the first argument (`arg_count`). This first argument can be accessed directly in the normal way—the `va_arg()` macro is used to get the others. This macro is passed the argument pointer, and the expected type of the argument. It evaluates to the current argument, and advances the argument pointer to the next one. In this case, the arguments are all `ints`, but they would be different types in an application like `printf()`. The `va_end()` macro tells the system that there are no more arguments to get (that the argument pointer has advanced to the end of the list).

The foregoing is implemented in `<stdarg.h>`, shown in Listing A.24. The argument-pointer type, `va_list`, is a pointer to `char`. It's declared as such because pointer arithmetic on a character pointer is just plain arithmetic. When you add 1 to a character pointer, you actually add 1 to the physical pointer (because the size of a character is 1). The best way to understand the other two macros is to watch them work. A typical C compiler passes arguments to subroutines by pushing them onto the run-time stack in reverse order. A call like

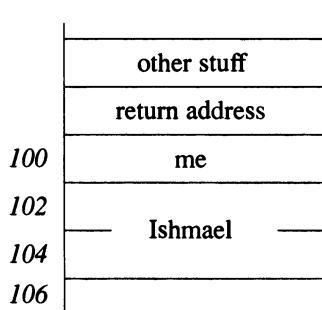
```

int me,
long Ishmael;

call( me, Ishmael );

```

looks like this on the stack (assuming a 16-bit `int` and a 32-bit `long`):



The `va_start(me, arg_ptr);` macro initializes `arg_ptr` to point at the second argument as follows: `&first` evaluates to the address of `me`—to 100—and `100+sizeof(first)` yields 102, the base address of the next argument on the stack.

Start up variable-argument processing, `va_start()`.

Get arguments, `va_arg()`.

End variable-argument processing, `va_end()`.

`va_list`, implemented as `char*`.

Implementing `va_start()`.

5. Note that the ellipsis in the `print_int()` argument list is not supported by UNIX C. The `VA_LIST` macro in `<debug.h>`, discussed on page 684, can be used to correct this deficiency. `VA_LIST` is used in various subroutines in subsequent sections.

Listing A.24. *stdarg.h*—Support for Variable-Argument Lists

```

1 typedef char *va_list;
2 #define va_start(arg_ptr,first) arg_ptr = (va_list)&first + sizeof(first)
3 #define va_arg(arg_ptr,type)    ((type *) (arg_ptr += sizeof(type)))[-1]
4 #define va_end(arg_ptr)        /* empty */

```

Implementing `va_arg()`. The cast is required in front of `&first` in order to defeat pointer arithmetic; otherwise you'd get 104, because `&first` evaluates to an `int` pointer by default. A subsequent `va_arg(arg_ptr,long)` call will fetch Ishmael and advance the pointer. It expands to:

```
((long *) (arg_ptr += sizeof(long)))[-1]
```

The `arg_ptr += sizeof(long)` yields 106 in this example. That is, it yields a pointer to just past the variable that you want. The `(long *) (106)` is treated as if it were a pointer to an array of longs that had been initialized to point at cell 106, and the `[-1]` gets the `long`-sized number in front of the pointer. You could think of the whole thing like this:

```

long      *p;
p = (long *)arg_ptr;           /* pointer to current argument */
++p;                         /* skip past it */
target = p[-1];              /* fetch it by backing up. */
arg_ptr += sizeof(long);     /* skip to next argument */

```

Note that some non-ANSI compilers can do the foregoing with:

```
* ((long *)arg_ptr) ++
```

but this statement isn't acceptable to many compilers (because the cast forms an rvalue, and `++` must be applied to an lvalue).

A.5 Conversion Functions

This section contains descriptions of several small data-conversion subroutines:

Convert string to long.

```
long          stol (char **str)
unsigned long stoul(char **str)
```

These routines are somewhat more powerful versions of the standard `atoi()`. `stol()` is passed the address of a character pointer (note the double indirection). It returns, in a `long`, the value of the number represented by the string, and updates the pointer to point past the number. If the string begins with a `0x`, the number is assumed to be hex; otherwise, if it begins with a `0`, the number is assumed to be octal; otherwise it is decimal. Conversion stops on encountering the first character that is not a digit in the indicated radix. Leading white space is ignored, and a leading minus sign is recognized. `stoul()` works much the same, except that it returns an `unsigned long`, and doesn't recognize a leading minus sign.

Translate escape sequences to binary.

```
int esc(char **s)
```

Returns the character associated with the “escape sequence” pointed to by `*s` and modifies `*s` to point past the sequence. The recognized strings are summarized in Table A.1.

Table A.1. Escape Sequences Recognized by `esc()`

Input string	Returned value	Notes
\b	0x08	Backspace.
\f	0x0c	Formfeed.
\n	0x0a	Newline.
\r	0x0d	Carriage return.
\s	0x20	Space.
\t	0x09	Horizontal tab.
\e	0x1b	ASCII ESC character.
\DDD	0x??	Number formed of 1–3 octal digits.
\xDDD	0x??	Number formed of 1–3 hex digits.
\^C	0x??	C is any letter. The equivalent control code is returned.
\c	c	a backslash followed by anything else returns the character following the backslash (and *s is advanced two characters).
c	c	Characters not preceded by a backslash are just returned (and *s is advanced 1 character).

```
char *bin_to_ascii(int c, int use_hex)
```

Convert binary to printing ASCII.

Return a pointer to a string that represents c in human-readable form. This string contains only the character itself for normal characters; it holds an escape sequence (\n, \t, \x00, and so forth) for others. A single quote ('') is returned as the two-character string "\'". The returned string is destroyed the next time `bin_to_ascii()` is called, so don't call it twice in a single `printf()` statement. If `use_hex` is true then escape sequences of the form \xDD are used for nonstandard control characters (D is a hex digit); otherwise, sequences of the form \DDD are used (D is an octal digit).

`stoul()` and `stol()` are implemented in Listing A.25. The `esc()` subroutine is in Listing A.26, and `bin_to_ascii()` is in Listing A.26 All are commented sufficiently that no additional description is needed here.

Listing A.25. *stol.c*—Convert String-to-Long

```

1 #include <ctype.h>
2
3 unsigned long stoul(instr)
4 char **instr;
5 {
6     /* Convert string to long. If string starts with 0x it is interpreted as
7      * a hex number, else if it starts with a 0 it is octal, else it is
8      * decimal. Conversion stops on encountering the first character which is
9      * not a digit in the indicated radix. *instr is updated to point past the
10     * end of the number.
11 */
12
13     unsigned long num = 0 ;
14     char *str = *instr;
15
16     while( isspace(*str) )
17         ++str ;
18 }
```



Listing A.25. continued...

```

19     if(*str != '0')
20     {
21         while( isdigit(*str) )
22             num = (num * 10) + (*str++ - '0');
23     }
24 else
25 {
26     if (++str == 'x' || *str == 'X')                  /* hex */
27     {
28         for( ++str; isxdigit(*str); ++str )
29             num = (num * 16) + ( isdigit(*str) ? *str - '0'
30                               : toupper(*str) - 'A' + 10 );
31     }
32 else
33 {
34     while( '0' <= *str && *str <= '7' )           /* octal */
35     {
36         num *= 8;
37         num += *str++ - '0' ;
38     }
39 }
40 }
41 *instr = str;
42 return( num );
43 }
44
45 /*-----*/
46
47 long    stol(instr)
48 char    **instr;
49 {
50     /* Like stoul(), but recognizes a leading minus sign and returns a signed
51      * long.
52      */
53
54     while( isspace(**instr) )
55         ++*instr ;
56
57     if( **instr != '-' )
58         return (long)( stoul(instr) );
59     else
60     {
61         ++*instr;                                /* Skip the minus sign. */
62         return -(long)( stoul(instr) );
63     }
64 }
```

Listing A.26. esc.c—Map Escape Sequences to Binary

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <tools/debug.h>
4
5 /* ESC.C      Map escape sequences to single characters */
6
```

Listing A.26. continued...

```

7  PRIVATE int hex2bin P(( int c ));
8  PRIVATE int oct2bin P(( int c ));
9  /*-----*/
10
11 #define ISHEXDIGIT(x) (isdigit(x) || ('a'<=(x)&&(x)<='f') || ('A'<=(x)&&(x)<='F'))
12 #define ISOCTDIGIT(x) ('0'<=(x) && (x)<='7')
13
14 PRIVATE int hex2bin(c)
15     int      c;
16 {
17     /* Convert the hex digit represented by 'c' to an int. 'c' must be one of
18      * the following characters: 0123456789abcdefABCDEF
19      */
20     return (isdigit(c) ? (c)-'0': ((toupper(c))-'A')+10) & 0xf;
21 }
22
23 PRIVATE int oct2bin(c)
24     int      c;
25 {
26     /* Convert the hex digit represented by 'c' to an int. 'c' must be a
27      * digit in the range '0'-'7'.
28      */
29     return ( ((c)-'0') & 0x7 );
30 }
31
32 /*-----*/
33
34 PUBLIC int      esc(s)
35     char     **s;
36 {
37     /* Map escape sequences into their equivalent symbols. Return the equivalent
38      * ASCII character. *s is advanced past the escape sequence. If no escape
39      * sequence is present, the current character is returned and the string
40      * is advanced by one. The following are recognized:
41      *
42      * \b      backspace
43      * \f      formfeed
44      * \n      newline
45      * \r      carriage return
46      * \s      space
47      * \t      tab
48      * \e      ASCII ESC character ('\033')
49      * \DDD    number formed of 1-3 octal digits
50      * \xDDD   number formed of 1-3 hex digits (two required)
51      * \^C     C = any letter. Control code
52      */
53
54     register int      rval;
55
56     if( **s != '\\\\' )
57         rval = *( (*s)++ );
58     else
59     {
60         ++(*s);                                /* Skip the \\ */
61         switch( toupper(**s) )
62         {
63             case '\\0': rval = '\\\\';           break;
64             case 'B':   rval = '\\b';            break;

```



Listing A.26. continued...

```

65      case 'F': rval = '\f' ; break;
66      case 'N': rval = '\n' ; break;
67      case 'R': rval = '\r' ; break;
68      case 'S': rval = ' ' ; break;
69      case 'T': rval = '\t' ; break;
70      case 'E': rval = '\033'; break;
71
72      case '^': rval = *++(*s) ;
73                  rval = toupper(rval) - '@' ;
74                  break;
75
76      case 'X': rval = 0;
77                  ++(*s);
78                  if( ISHEXDIGIT(**s) )
79                  {
80                      rval = hex2bin( *(s)++ );
81                  }
82                  if( ISHEXDIGIT(**s) )
83                  {
84                      rval <= 4;
85                      rval |= hex2bin( *(s)++ );
86                  }
87                  if( ISHEXDIGIT(**s) )
88                  {
89                      rval <= 4;
90                      rval |= hex2bin( *(s)++ );
91                  }
92                  --(*s);
93                  break;
94
95      default: if( !ISOCTDIGIT(**s) )
96                  rval = **s;
97                  else
98                  {
99                      ++(*s);
100                     rval = oct2bin( *(s)++ );
101                     if( ISOCTDIGIT(**s) )
102                     {
103                         rval <= 3;
104                         rval |= oct2bin( *(s)++ );
105                     }
106                     if( ISOCTDIGIT(**s) )
107                     {
108                         rval <= 3;
109                         rval |= oct2bin( *(s)++ );
110                     }
111                     --(*s);
112                 }
113                 break;
114             }
115             ++(*s);
116         }
117         return rval;
118     }

```

Listing A.27. *bintoasc.c*— Convert Binary to Human-Readable String

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/compiler.h>      /* for prototypes only */
4
5 char *bin_to_ascii( c, use_hex )
6 {
7     /* Return a pointer to a string that represents c. This will be the
8      * character itself for normal characters and an escape sequence (\n, \t,
9      * \x00, etc., for most others). A ' is represented as '\'. The string will
10     * be destroyed the next time bin_to_ascii() is called. If "use_hex" is true
11     * then \xDD escape sequences are used. Otherwise, octal sequences (\DDD)
12     * are used. (see also: pchar.c)
13     */
14
15     static unsigned char buf[8];
16
17     c &= 0xff ;
18     if( ' ' <= c && c < 0x7f && c != '\\' && c != '\\\\')
19     {
20         buf[0] = c;
21         buf[1] = '\\0';
22     }
23     else
24     {
25         buf[0] = '\\\\' ;
26         buf[2] = '\\0' ;
27
28         switch(c)
29         {
30             case '\\\\': buf[1] = '\\\\'; break;
31             case '\\\'': buf[1] = '\\\''; break;
32             case '\\b': buf[1] = 'b' ; break;
33             case '\\f': buf[1] = 'f' ; break;
34             case '\\t': buf[1] = 't' ; break;
35             case '\\r': buf[1] = 'r' ; break;
36             case '\\n': buf[1] = 'n' ; break;
37             default : sprintf( &buf[1], use_hex ? "x%03x" : "%03o", c );
38                         break;
39         }
40     }
41     return buf;
42 }
```

A.6 Print Functions

This section describes several output functions, many of which use the ANSI variable-argument mechanism and the conversion functions described in the last two sections. They are organized functionally.

void ferr(char *fmt, ...)

Print error message and exit.

This routine is a version of `printf()` for fatal error messages. It sends output to `stderr` rather than `stdout`, and it doesn't return; rather, it terminates the program with the following call:

```
exit( on_ferr() );
```

(`on_ferr()` is described below). Normally, the routine works like `printf()` in other respects. If, however, it is called with `ferr(NULL, string)`, it works like `perror(string)` in that it prints an error message associated with the most recent I/O-system error. The program still terminates as described earlier, however. The source code for `ferr()` is in Listing A.28. The `prnt()` call on line 19 will be described in a moment.

Listing A.28. *ferr.c*— Fatal-Error Processing

```

1 #include <stdio.h>
2 #include <stdlib.h>      /* errno is declared here      */
3 #include <stdarg.h>
4 #include <tools/debug.h> /* VA_LIST definition is here */
5 #include <tools/l.h>     /* Needed only for prototype */
6
7 /* Note that ferr() is typed as int, even though it usually doesn't return,
8 * because it's occasionally useful to use it inside a conditional expression
9 * where a type will be required. VA_LIST expands to ... if ANSI is #defined,
10 * otherwise it expands to _a_r_g_s_.
11 */
12
13 int      ferr( fmt, VA_LIST )
14 char    *fmt;
15 {
16     va_list      args;
17     va_start( args, fmt );
18
19     if( fmt ) prnt   ( fputc, stderr, fmt, args );
20     else      perror ( va_arg(args, char* ) );
21
22     va_end( args );
23     exit( on_ferr() );
24 }
```

Error handler for `ferr()`. **int on_ferr(void)**

This routine is the default error handler called by `ferr()` just before exiting. It returns the current contents of the system `errno` variable, which is, in turn, passed back up to the operating system as the process' exit status. You can provide your own `on_ferr()` to preempt the library version. The source code is in Listing A.29.

Print human-readable
string.

void fputstr(char *str, int maxlen, FILE *stream)

This function writes a string (`str`) having at most `maxlen` characters to the indicated `stream` in human-readable form. All control characters, and so forth, are mapped to printable strings using the `bin_to_ascii()` conversion function described earlier (on page 727). The source code is in Listing A.30.

Print human-readable
character.

void pchar(FILE *stream, int c)

This function works like `putc()`, except that control characters are mapped to human-readable strings using `bin_to_ascii()`. The source code is in Listing A.31.

Listing A.29. *onferr.c*—Action Function for ferr()

```

1 #include <stdlib.h>
2 #include <tools/debug.h>
3 #include <tools/l.h> /* Needed only for prototypes */
4
5 /* This is the default routine called by ferr when it exits. It should return
6 * the exit status. You can supply your own version of this routine if you like.
7 */
8 int on_ferr()
9 {
10     extern int errno;
11     return errno;
12 }
```

Listing A.30. *fputstr.c*—Fatal-Error Processing—Service Routine

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/compiler.h>
4
5 /* FPUTSTR.C: Print a string with control characters mapped to readable strings.
6 */
7
8 void fputstr( str, maxlen, stream )
9 char *str;
10 FILE *stream;
11 {
12     char *s;
13
14     while( *str && maxlen >= 0 )
15     {
16         s = bin_to_ascii( *str++, 1 );
17         while( *s && --maxlen >= 0 )
18             putc( *s++ , stream );
19     }
20 }
```

Listing A.31. *pchar.c*—Print Character in Human-Readable Form

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/compiler.h>
4
5 void pchar(c, stream)
6 FILE *stream;
7 {
8     fputs( bin_to_ascii( c, 1 ), stream );
9 }
```

void printv(FILE *stream, **char** **argv)

Print vector array.

This function prints an argv-like array of pointers to strings to the indicated stream, one string per line (the '\n' is inserted automatically at the end of every string). The source code is in Listing A.32.

Listing A.32. *printv.c*—Print argv-like Vector Array

```

1 #include <stdio.h>
2
3 void printv( fp, argv )
4 FILE *fp;
5 char **argv;
6 {
7     /* Print an argv-like array of pointers to strings, one string per line.
8      * The array must be NULL terminated.
9      */
10    while( *argv )
11        fprintf( fp, "%s\n", *argv++ );
12 }
13
14 void comment( fp, argv )
15 FILE *fp;
16 char **argv;
17 {
18     /* Works like printv except that the array is printed as a C comment. */
19
20     fprintf(fp, "\n/*-----\n");
21     while( *argv )
22         fprintf(fp, " * %s\n", *argv++ );
23     fprintf(fp, " */\n\n");
24 }
```

Print multiple-line comment.

void comment(**FILE** *stream, **char** **argv)

This function also prints an argv-like array of pointers to strings to the indicated stream, one string per line. The output text is put into a C comment, however. Output takes the following form:

```

/*-----
 * string in argv[0]
 * string in argv[1]
 * ....
 * string in argv[N]
 */
```

The source code is also in Listing A.32.

A printf() workhorse function. **void** prnt(**int** (*ofunct)(), **void** *ofunct_arg, **char** *format, **va_list** args)

prnt() is a variant on the UNIX_doprnt() and ANSI vfprintf() functions that lets you write printf()-like output routines in a portable way. It is passed a pointer to a single-character output function (ofunct), a parameter that is relayed to this output function (ofunct_arg), a pointer to a format string, and a pointer to the location on the run-time stack where the other arguments are found. This last parameter is usually derived using the va_start() macro in <stdlib.h>, described earlier.

To see how prnt() is used, fprintf() can be implemented as follows:

```

#include <stdio.h>
#include <stdarg.h>

fprintf( stream, format, ... );
FILE      *stream;
char     *format;
{
    extern int fputc();
    va_list    args;
    va_start (args, format );
    prnt     (fputc, stream, format, args );
    va_end   (args);
}

```

sprintf() can be implemented as follows:

```

putstr(c, p)
int      c;
char    **p;
{
    * (p)++ = c ;
}
sprintf(str, format, ...)
char    *str, *format;
{
    va_list    args;
    va_start (args, format );
    prnt     (putstr, &str, format, args );
    *str = '\0' ;
    va_end   (args);
}

```

The prnt() subroutine is required by the curses implementation discussed below, which must be able to change the output subroutine at will. Neither the ANSI nor the UNIX function gives you this capability. The source code for prnt() is described in depth below.

void stop_prnt(**void**)

Clean up after prnt().

This routine must be called by all programs that use prnt() just before termination [after the last prnt() call].

The prnt(), subroutine is in Listing A.33. Two versions are presented, one for the ANSI environment and another for UNIX. The first version (on lines 19 to 31) produces an ANSI-compatible function. It uses vsprintf() to do the conversion into a buffer, and then prints the buffer. Though this version is desirable in that it supports all the printf() conversions, it has a drawback. The buffer requires a modicum of stack space, and the Microsoft implementation of vsprintf() also uses a lot of run-time stack. The second, UNIX compatible version of prnt() is on lines 37 to 65 of Listing A.33. There are two problems here: First, the UNIX variable-argument mechanism is different from the ANSI one, so slightly different procedures must be used to get arguments off the stack. Second, there is no UNIX equivalent to vsprintf(). The only available printf driver is _doprnt(), which works like the ANSI vprintf() function except that the arguments are in different positions. In order to use an arbitrary output function, you must format to a FILE, rewind the file, and then transfer characters one at a time to the desired output function. Ugh,⁶ you may well say. This approach is

ANSI version of prnt().

UNIX version of prnt().

6. That's a technical term.

actually not as much of a kludge as it seems because the number of characters written and read is so small. More often than not, there's no disk activity at all because the reads and writes won't get past the I/O system's internal buffering.

Listing A.33. prnt.c—General-Purpose printf() Driver

```

1 #include <stdio.h>
2 #include <tools/debug.h>
3 #include <tools/l.h>
4
5 /-----
6 * Glue formatting workhorse functions to various environments. One of three
7 * versions of the workhorse function is used, depending on various #defines:
8 *
9 * if ANSI is defined      vsprintf()  Standard ANSI function
10 * if ANSI is not defined _doprnt()  Standard UNIX function
11 *
12 * The default with Microsoft C is MSDOS defined and ANSI not defined,
13 * so idoprnt() will be used unless you change things with explicit macro
14 * definitions.
15 */
16 #ifdef ANSI /*-----*/
17 #include <stdarg.h>
18
19 PUBLIC void prnt( ofunct, funct_arg, format, args )
20 int     (*ofunct)();
21 void    *funct_arg;
22 char   *format;
23 va_list args;
24 {
25     char buf[256], *p ;
26     int  vsprintf( char* buf, char* fmt, va_list args );
27
28     vsprintf(buf, format, args);
29
30     for( p = buf; *p ; p++ )
31         (*ofunct)( *p, funct_arg );
32 }
33
34 PUBLIC void stop_prnt(){}
35
36 #else /* UNIX -----*/
37 #include <varargs.h>
38
39 static FILE    *Tmp_file = NULL ;
40 static char   *Tmp_name ;
41
42 PUBLIC void prnt( ofunct, funct_arg, fmt, argp )
43 int     (*ofunct)();
44 void    *funct_arg;
45 char   *fmt;
46 int     *argp;
47 {
48     int   c;
49     char  *mktemp();
50
51     if( !Tmp_file )
52         if( !(Tmp_file = fopen( Tmp_name = mktemp("yyXXXXXX"), "w+" ) ) )
53         {
54             fprintf(stderr,"Can't open temporary file %s\n", Tmp_name );
55             exit( 1 );

```



Listing A.33. continued...

```

56         }
57
58     _doprnt( fmt, argp, Tmp_file );
59     putc( 0, Tmp_file );
60     rewind( Tmp_file );
61
62     while( (c = getc(Tmp_file)) != EOF && c )
63         (*ofunct)( c, funct_arg );
64     rewind( Tmp_file );
65 }
66
67 PUBLIC void stop_prnt()
68 {
69     fclose( Tmp_file );           /* Remove prnt temporary file */
70     unlink( Tmp_name );
71     Tmp_file = NULL;
72 }
73
74 /*-----*/
75
76 PUBLIC void vfprintf( stream, fmt, argp )
77 FILE *stream;
78 char *fmt, *argp;
79 {
80     _doprnt( fmt, argp, stream );
81 }
82
83 PUBLIC void vprintf( fmt, argp )
84 char *fmt, *argp;
85 {
86     _doprnt( fmt, argp, stdout );
87 }
88
89 PRIVATE void putstr(c, p)
90 int c;
91 char **p;
92 {
93     *( *p )++ = c ;
94 }
95
96 PUBLIC void vsprintf( str, fmt, argp )
97 char *str, *fmt, *argp;
98 {
99     prnt( putstr, &str, fmt, argp );
100    *str = '\0' ;
101 }
102 #endif

```

The temporary file is created automatically on lines 51 to 56 of Listing A.34 the first time `prnt()` is called. You must call `stop_prnt()` (on lines 67 to 72) to delete the file. You can do so after every `prnt()` call, but it's much more efficient to do it only once, just before the program terminates. You'll lose a `FILE` pointer in this case, but it saves a lot of time.

The final part of Listing A.33 (lines 76 to 101) contains UNIX versions of the ANSI `vprintf()`, `vfprintf()`, and `vsprintf()` functions.

`ssort()` is like
`qsort()`.

Sorting `argv` with
`ssort()`.

Sorting an array of structures with `ssort()`.

A.7 Sorting

This sections describes two functions for sorting in-memory arrays: `ssort()` and `assort()`. Both are modeled after the UNIX `qsort()` function. I've provided alternate sorting functions because most `qsort()` implementations use a "quick sort", which does not behave particularly well with small arrays or arrays that might already be sorted. A Shell sort, which does not have these problems, is used here. Both sort functions are general-purpose functions that can be used on arrays of any type of object—the same sorting function can sort arrays of `ints`, arrays of character pointers, arrays of structures, and so forth. Taking `ssort()` as characteristic, `argv` is sorted as follows:

```
cmp( p1, p2 )
char **p1, **p2;
{
    return strcmp( *p1, *p2 );
}

main( argc, argv )
int argc;
char **argv;
{
    ssort( argv, argc, sizeof(*argv), cmp );
}
```

The sort function is passed the base address of the array to be sorted, the number of elements in the array, the size of one element, and a pointer to a comparison function. This comparison function is passed pointers to two array elements, and it should otherwise work like `strcmp()`, returning a negative number if the key associated with the first argument is less than the key associated with the second, zero if the two keys are equal, and a positive number if the second key is larger. You can use `ssort()` to sort an array of structures as follows:

```
typedef struct
{
    char *key;
    int other_stuff;
} record;

sort_cmp( p1, p2 )
record *p1, *p2;
{
    return( strcmp(p1->key, p2->key) );
}

plato()
{
    record field[ 10 ];
    ...
    ssort( field, 10, sizeof(record), sort_cmp );
}
```

Of course, it's usually better to sort arrays of pointers to structures rather than arrays of structures—there's a lot less to move when two array elements are swapped. The calling conventions for the two sort routines are as follows:

Sort array.

`void ssort(void *base, int nel, int elesize, int (*cmp)())`

Sort an array at base address `base`, having `nel` elements, each of `elesize` bytes. The comparison function, `cmp`, is passed pointers to two array elements and should return as follows:

cmp(p1, p2);	
*p1 < *p2	return a negative number
*p1 = *p2	return zero
*p1 > *p2	return a positive number

The `ssort()` function is identical in calling syntax to the standard `qsort()` function. A Shell sort is used by `ssort()`, however, rather than a quicksort. The Shell sort is more appropriate for use on arrays with a small number of elements and arrays that might already be sorted. Also, since Shell sort is nonrecursive, it is a safer function to use than quicksort, which can cause stack overflows at run time.

```
void assort(void **base, int nel, int elesize, int (*cmp) ())      Sort array of pointers.
```

This routine is a version of `ssort()` optimized to sort arrays of pointers. It takes the same arguments as `ssort()` so that the routines can be used interchangeably on arrays of pointers. The `base` argument to `assort()` must reference an array of pointers, however, and the `elesize` argument is ignored.

A.7.1 Shell Sort—Theory

The sorting method used by `ssort()` is a Shell sort, named after its inventor, Donald Shell. It is essentially an improved bubble sort, so I'll digress for a moment and describe bubble sort. A bubble-sort that arranges an array of `ints` into ascending order looks like this:

Bubble sort.

```
int array[ ASIZE ];
int i, j, temp ;

for( i = 1; i < ASIZE; ++i )
    for( j = i-1; j >= 0; --j )
        if( array[j] > array[j+1] )
            swap( array+i, array+j ); /* swap array[i] and array[j] */
```

The outer loop controls the effective array size. It starts out sorting a two-element array, then it sorts a three-element array, and so forth. The inner loop moves the new element (the one added when the effective array size was increased) to its correct place in the previously-sorted array. Consider a worst case sort (where the array is already sorted, but in reverse order):

5 4 3 2 1

In the first pass, the outer loop starts out with a two-element array:

5 4 3 2 1

And the inner loop swaps these two elements because they're out of place:

4 5 3 2 1

The next pass increases the effective array size to three elements:

4 5 3 2 1

And the 3 is moved into its proper place like this:

```
4 5 3                2 1
4 3 5                2 1
3 4 5                2 1
```

The rest of the sort looks like this:

3 4 5 2	1
3 4 2 5	1
3 2 4 5	1
2 3 4 5	1
2 3 4 5 1	
2 3 4 1 5	
2 3 1 4 5	
2 1 3 4 5	
1 2 3 4 5	

As you can see, this worst-case sort is very inefficient. In an N -element array, roughly N^2 swaps (and as many comparisons) are required to get the array sorted. Even the average case requires N^2 comparisons, even though it will use fewer swaps. This behavior is quite measurable in most computer programs, and slows down the program unnecessarily.

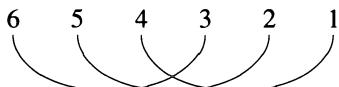
Shell sort.

The Shell sort improves the bubble sort by trying to move the most out-of-place elements into the proper place as quickly as possible, instead of letting the out-of-place element percolate through the array one place at a time. Other sort strategies (such as the quicksort used for most `qsort()` implementations) are theoretically more efficient, but the overhead required by these other methods is often great enough that the theoretically less-efficient Shell sort is faster (at least with small arrays—less than a hundred elements or so).

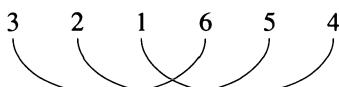
The basic strategy of Shell sort is to partition the array into several smaller subarrays whose elements are spread out over the original array. The subarrays are sorted using a bubble sort, and then the array is repartitioned into a smaller number of subarrays, each having more elements. The process is continued until there's only one subarray comprising the entire array. For example, if the initial array looks like this:

6 5 4 3 2 1

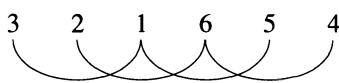
You can partition it into three, two-element arrays like this:



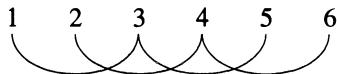
The first array is [6,3], the second is [5,2] and the third holds [4,1]. The distance between any two elements in the subarray is called the *gap size*. Here, the gap size is 3. The three subarrays are now sorted. The 6 and 3 will be swapped, as will the 5 and 2, and the 4 and 1, yielding the following:



The array is repartitioned, this time with a gap size of two:



There are now two, three-element arrays having the members [3, 1, 5] and [2, 6, 4]. These are sorted, swapping the 1 and 3, and the 4 and 6, yielding.



The gap size is now reduced to 1, yielding:



but, since this array is already sorted, nothing more needs to be done. Note that only 5 swaps were used here rather than the 15 swaps that would be required by the bubble sort. In the average case, roughly $N^{1.2}$ swaps are required to sort an array of size N .⁷

A.7.2 Shell Sort—Implementation

The `ssort()` function in Listing A.34 is a general-purpose Shell-sort function. The initial gap size is selected on line 16 from a number in the series: 1, 4, 13, 40, 121, ... ($(3N+1)$). The largest number in this series that is less than or equal to the array size is used. (This is the series recommended by Knuth in *The Art of Computer Programming* as an optimal choice. Note that an even power of two, as is used in many Shell-sort implementations is among the worst choice of gap sizes.) The `for` statement on line 19 controls the size of the subarrays. The gap is divided by three on each pass, yielding the previous element in the foregoing series ($40/3=13$, $13/3=4$, $4/3=1$ —I'm using integer division). The two loops on lines 20 to 33 are doing a bubble-sort pass on the partitioned subarray.

Choosing the gap size.

Since the compiler doesn't know the size of an array element at compile time, it can't do the pointer arithmetic; we have to do the pointer arithmetic ourselves by declaring the array to be a character array on line eight and then explicitly multiplying by the size of an array element every time a pointer is advanced (on lines 23 and 24). The array elements also have to be swapped one byte at a time. The swap time could be improved by passing a pointer to a swap function as well as a comparison function to `ssort`, but I wanted to maintain compatibility with the standard `qsort()`, and so didn't change any of the parameters.

Doing explicit pointer arithmetic.

One change that I did make is shown in Listing A.35. `asort()` is a version of `ssort()` that sorts only arrays of pointer-sized objects—the vast majority of cases. The routine takes the same arguments as `ssort()`, but it ignores `elsize`.

Implementing `assort()`.

A.8 Miscellaneous Functions

This section discusses six routines that don't fit nicely into any of the other categories:

```
int copyfile(char *dst, char *src, char *mode)
int movefile(char *dst, char *src, char *mode)
```

Copy entire file.
Move entire file.

`copyfile()` copies the contents of the file named in `src` to the file named in

7. See [Knuth], vol. 3, pp. 84f.

Listing A.34. *ssort.c*—General-Purpose Shell Sort

```

1  /* SSORT.C           Works just like qsort() except that a shell sort, rather
2   * than a quick sort, is used. This is more efficient than
3   * quicksort for small numbers of elements, and it's not recursive (so will use
4   * much less stack space).
5   */
6
7  void      ssort( base, nel, elsize, cmp )
8  char     *base;
9  int       nel, elsize;
10 int      (*cmp)();
11 {
12     int      i, j;
13     int      gap, k, tmp ;
14     char    *p1, *p2;
15
16     for( gap=1; gap <= nel; gap = 3*gap + 1 )
17         ;
18
19     for( gap /= 3; gap > 0 ; gap /= 3 )
20         for( i = gap; i < nel; i++ )
21             for( j = i-gap; j >= 0 ; j -= gap )
22             {
23                 p1 = base + ( j      * elsize);
24                 p2 = base + ((j+gap) * elsize);
25
26                 if( (*cmp)( p1, p2 ) <= 0 )          /* Compare two elements */
27                     break;
28
29                 for( k = elsize; --k >= 0 ; )        /* Swap two elements, one */
30                     {                                /* byte at a time. */
31                         tmp    = *p1;
32                         *p1++ = *p2;
33                         *p2++ = tmp;
34                     }
35             }
36 }
```

dst. If the mode argument is "w", the destination is overwritten when it already exists; otherwise, the contents of the source file are appended to the end of the destination. The return values are as follows:

- 0 Copy was successful.
- 1 Destination file couldn't be opened.
- 2 Source file couldn't be opened.
- 3 Read error while copying.
- 4 Write error while copying.

`movefile()` works like `copyfile()`, except that the source file is deleted if the copy is successful. `movefile()` differs from the standard `rename()` function in that it allows a move across devices (from one disk to another), and it supports an append mode—`rename()` is faster if you're moving a file somewhere else on the same device, and `rename()` can be used on directory names, however. The sources for `copyfile()` and `movefile()` are in Listings A.36 and A.37.

Listing A.35. assort.c—ssort() Optimized for Arrays of Pointers

```

1  /* ASSORT.C          A version of ssort optimized for arrays of pointers. */
2
3  void    assort( base, nel, elsize, cmp )
4  void    **base;
5  int     nel;
6  int     elsize;      /* ignored */
7  int     (*cmp) ();
8  {
9      int     i, j, gap;
10     void   *tmp, *p1, **p2;
11
12     for( gap=1; gap <= nel; gap = 3*gap + 1 )
13         ;
14
15     for( gap /= 3; gap > 0 ; gap /= 3 )
16         for( i = gap; i < nel; i++ )
17             for( j = i-gap; j >= 0 ; j -= gap )
18             {
19                 p1 = base + ( j );
20                 p2 = base + ((j+gap) );
21
22                 if( (*cmp)( p1, p2 ) <= 0 )
23                     break;
24
25                 tmp = *p1;
26                 *p1 = *p2;
27                 *p2 = tmp;
28             }
29 }
```

Listing A.36. copyfile.c—Copy Contents of File

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <errno.h>
4
5 #include <sys/types.h>
6 #include <sys/stat.h>
7 #include <io.h>
8
9 #define ERR_NONE          0
10 #define ERR_DST_OPEN      -1
11 #define ERR_SRC_OPEN      -2
12 #define ERR_READ          -3
13 #define ERR_WRITE         -4
14
15 copyfile( dst, src, mode )
16 char   *dst, *src;
17 char   *mode;           /* "w" or "a" */
18 {
19     /* Copy the src to the destination file, opening the destination in the
20      * indicated mode. Note that the buffer size used on the first call is
21      * used on subsequent calls as well. Return values are defined, above.
22      * errno will hold the appropriate error code if the return value is <0.
23      */
24 }
```



Listing A.36. continued...

```

25     int          fd_dst, fd_src;
26     char         *buf;
27     int          got;
28     int          werror;
29     int          ret_val = ERR_NONE ;
30     static unsigned size    = 31 * 1024 ;
31
32     while( size > 0   &&  !(buf = malloc(size)) )
33         size -= 1024;
34
35     if( !size )           /* couldn't get a buffer, do it one byte at a time */
36     {
37         size = 1;
38         buf = "xx";      /* allocate a buffer implicitly */
39     }
40
41     fd_src = open(src, O_RDONLY | O_BINARY );
42     fd_dst = open(dst, O_WRONLY | O_BINARY | O_CREAT |
43                           (*mode=='w' ? O_TRUNC : O_APPEND),
44                           S_IREAD|S_IWRITE);
45
46     if      ( fd_src == -1 ){ ret_val = ERR_SRC_OPEN; }
47     else if( fd_dst == -1 ){ ret_val = ERR_DST_OPEN; }
48     else
49     {
50         while( (got = read(fd_src, buf, size)) > 0 )
51             if( (werror = write(fd_dst, buf, got)) == -1 )
52             {
53                 ret_val = ERR_WRITE;
54                 break;
55             }
56
57         if( got == -1 )
58             ret_val = ERR_READ;
59     }
60
61     if( fd_dst != -1 ) close ( fd_dst );
62     if( fd_src != -1 ) close ( fd_src );
63     if( size > 1    ) free   ( buf    );
64
65     return ret_val;
66 }
```

Listing A.37. movefile.c—Move File to Different Device

```

1 movefile( dst, src, mode )      /* Works like copyfile() (see copyfile.c)   */
2     char      *dst, *src;        /* but deletes src if the copy is successful */
3     char      *mode;
4 {
5     int    rval;
6     if( (rval = copyfile(dst, src, mode)) == 0 )
7         unlink( src );
8     return rval;
9 }
```

int *memiset(**int** *dst, **int** with_what, **int** count)

Fill memory with integer values.

This function is a version of the standard `memset()` function but it works on integer-sized objects rather than bytes. It fills `count` `ints`, based at `dst`, with the pattern in `with_what` and returns `dst`. The source code is in Listing A.38.

Listing A.38. *memiset.c*— Initialize Array of **int** to Arbitrary Value

```

1  /* Works like memset but fills integer arrays with an integer value      */
2  /* The count is the number of ints (not the number of bytes).                */
3
4  int      *memiset( dst, with_what, count )                                */
5  int      *dst, with_what, count;
6  {
7      int *targ;
8      for( targ = dst; --count >= 0 ; *targ++ = with_what )
9      ;
10     return dst;
11 }
```

int concat(**int** size, **char** *dst, ...)

Concatenate strings.

The `concat()` function concatenates an arbitrary number of strings into a single destination array (`dst`) of the indicated `size`. At most `size-1` characters are copied. All arguments following the `dst` argument are the source strings to be concatenated, and the list should end with a `NULL`. For example, the following code loads the `english` array with the string `angles, saxons, jutes`:

```

#include <stdio.h>    /* For NULL definition */

char target[ SIZE ];

concat( SIZE, target, "angles, ", "saxons, ", " jutes", NULL );
```

The second and third arguments can be the same, but the target-string pointer cannot appear in any of the other arguments. The following concatenates `new_string` to the end of `target`. This usage is easier to use than the standard `strncat()` function in many situations because it doesn't require you to keep a tally of the unused space in the target array.

```
concat( SIZE, target, target, new_string );
```

The source code is in Listing A.39. The amount of available space is returned, or `-1` if the string was truncated.

void searchenv(**char** *filename, **char** *env_name, **char** *pathname)

Search for file along directory path listed in an environment string.

This function searches for a specific file (`filename`), first in the current directory, and then along a path specified in the environment variable whose name is in `env_name`. The variable should hold a semicolon or space-delimited list of directory names. If the file is found, the full path name (including the file-name component) is put into `pathname`; otherwise, `pathname` will contain an empty, null-terminated string. The source code is in Listing A.40.

FILE *driver_1(**FILE** *output, **int** lines, **char** *fname)

Copy driver-template file.

int driver_2(**FILE** *output, lines)

These routines work together to transfer a template file to a L^EX or `occs` output file. `driver_1()` must be called first. It searches for the file named `fname`

Listing A.39. concat.c—Concatenate Strings

```

1 #include <stdio.h>
2 #include <stdarg.h>
3 #include <debug.h>      /* VA_LIST definition */
4
5 int     concat( size, dst, VA_LIST )
6 int     size;
7 char    *dst;
8 {
9     /* This subroutine concatenates an arbitrary number of strings into a single
10    * destination array (dst) of size "size." At most size-1 characters are
11    * copied. Use it like this:
12    *         char target[SIZE];
13    *         concat( SIZE, target, "first ", "second ",..., "last", NULL);
14    */
15
16    char    *src;
17    va_list args;
18    va_start( args, dst );
19
20    while( (src = va_arg(args,char *)) ) && size > 1 )
21        while( *src && size-- > 1 )
22            *dst++ = *src++ ;
23
24    *dst++ = '\0';
25    va_end( args );
26    return (size <= 1 && src && *src) ? -1 : size ;
27 }
```

by looking first in the current directory and then in any directory on the path specified by the LIB environment. This environment can list several, semicolon-delimited directory names.

Ctrl-L delimits parts of driver-template file.

The file should contain one or more Ctrl-L-delimited parts. The first part is copied by driver_1() to the stream indicated by output. If lines is true, then a #line directive that references the template file's current line number is output just before the block is output. NULL is returned if the template file can't be opened, otherwise the FILE pointer for the template file is returned. You can use this pointer to close the file after it has been copied.

All other Ctrl-L-delimited parts of the template file are printed by successive calls to driver_2(). One part is printed every time it's called. 1 is returned normally, 0 at end of file (if there are no more parts to print).

Use @ to mark comments in driver-template file.

Lines that begin with an @ sign are ignored by both subroutines. It's a fatal error to call driver_2() without a previously successful driver_1() call.

The source code for both routines is in Listing A.41.

A.9 Low-Level Video I/O Functions for the IBM PC

This section presents a set of very low-level terminal I/O functions for the IBM PC. They are used by the *curses* implementation presented in the next section. The routines in this section are the only ones in the *curses* package that are system dependent. As a consequence, I've not bothered to make them portable, because they'll always have to be rewritten if you port the code to another compiler.

Listing A.40. *searchen.c*— Search for File Along Path Specified in Environment

```

1  #include <stdio.h>
2
3  #define PBUF_SIZE 129           /* Maximum length of a path name - 1 */
4
5  void    searchenv( filename, envname, pathname )
6  char    *filename;          /* file name to search for */                      */
7  char    *envname;          /* environment name to use as PATH */                  */
8  char    *pathname;         /* Place to put full path name when found */      */
9  {
10     /* Search for file by looking in the directories listed in the envname
11     * environment. Put the full path name (if you find it) into pathname.
12     * Otherwise set *pathname to 0. Unlike the DOS PATH command (and the
13     * microsoft _searchenv), you can use either a space or semicolon
14     * to separate directory names. The pathname array must be at least
15     * 128 characters.
16     */
17
18     char pbuf[PBUF_SIZE];
19     char *p ;
20     char *strpbrk(), *strtok(), *getenv();
21
22     strcpy( pathname, filename );
23     if( access( pathname, 0 ) != -1 )           /* check current directory */
24         return;                                /* ...it's there. */
25
26     /* The file doesn't exist in the current directory. If a specific path was
27     * requested (ie. file contains \ or /) or if the environment isn't set,
28     * return a NULL, else search for the file on the path.
29     */
30
31     if( strpbrk(filename,"\\/") || ! (p = getenv(envname)) )
32     {
33         *pathname = '\0';
34         return;
35     }
36
37     strncpy( pbuf, p, PBUF_SIZE );
38     if( p = strtok( pbuf, ";" ) )
39     {
40         do
41         {
42             sprintf( pathname, "%0.90s\\%0.20s", p, filename );
43
44             if( access( pathname, 0 ) >= 0 )           /* found it */
45                 return;
46         }
47         while( p = strtok( NULL, ";" ) );
48     }
49     *pathname = '\0' ;
50 }
```

Two sets of complementary routines are presented here: a set of direct-video functions that write directly into the IBM's display memory (they are MGA and CGA compatible), and a set of similar routines that use the video functions built into the IBM-PC ROM-BIOS to do the I/O. The direct-video routines are faster, the BIOS routines are, at least in theory, more portable. (The direct-video functions work better than the BIOS ones in clones that have nonstandard BIOS implementations.)

Listing A.41. driver.c— Copy Driver Template to Output File

```

1 #include <stdio.h>
2 #include <ctype.h>
3 #include <stdlib.h>
4 #include <tools/debug.h>
5 #include <tools/compiler.h>      /* for prototypes */
6
7 /*-----*/
8
9 PUBLIC FILE *driver_1( FILE *output, int line, char *file_name );
10 PUBLIC int driver_2( FILE *output, int line );
11
12 PRIVATE FILE *Input_file = NULL ;
13 PRIVATE int Input_line;           /* line number of most-recently read line */
14 PRIVATE char File_name[80];       /* template-file name */
15
16 /*-----*/
17
18 PUBLIC FILE *driver_1( output, lines, file_name )
19 FILE    *output;
20 char   *file_name;
21 {
22     char path[80];
23
24     if( !(Input_file = fopen( file_name, "r" )) )
25     {
26         searchenv( file_name, "LIB", path );
27         if( !path || !(Input_file = fopen( path, "r")) )
28             return NULL;
29     }
30
31     strncpy( File_name, file_name, sizeof(File_name) );
32     Input_line = 0;
33     driver_2( output, lines );
34     return Input_file;
35 }
36
37 /*-----*/
38
39 PUBLIC int driver_2( output, lines )
40 FILE    *output;
41 {
42     static char buf[ 256 ];
43     char      *p;
44     int       processing_comment = 0;
45
46     if( !Input_file )
47         ferr( "INTERNAL ERROR [driver_2], Template file not open.\n" );
48
49     if( lines )
50         fprintf( output, "\n#line %d \"%s\"\n", Input_line + 1, File_name );
51
52     while( fgets(buf, sizeof(buf), Input_file) )
53     {
54         ++Input_line;
55         if( *buf == '\f' )
56             break;
57

```



Listing A.41. continued...

```

58         for( p = buf; isspace(*p); ++p )
59             ;
60             if( *p == '@' )
61             {
62                 processing_comment = 1;
63                 continue;
64             }
65             else if( processing_comment )           /* Previous line was a comment, */
66             {                                     /* but current line is not.      */
67                 processing_comment = 0;
68                 if( lines )
69                     fprintf( output, "\n#line %d \"%s\"\n", Input_line, File_name );
70                 }
71             fputs( buf, output );
72         }
73     return( feof(Input_file) );
74 }
```

A.9.1 IBM Video I/O—Overview

You need to know a little about low-level I/O to understand the function descriptions that follow. The IBM PC uses several different hardware adapters for video output. The most common are the Color Graphics Adapter (CGA) and Monochrome Graphics Adapter (MGA). Most other interface cards can emulate one or the other of these. Both the MGA and CGA use a block of memory to represent the screen—a two-dimensional array in high memory holds the characters that are currently being displayed. The position of the character in the array determines the position of the character on the screen. Several such arrays, called *video display pages*, are available, but only one page can be displayed at a time. Both systems use a 25×80 array, but the MGA bases the default video page at absolute memory location 0xb0000 (B000:0000), and the CGA puts it at 0xb8000. The other pages are at proportionally higher addresses.

Both cards use a 16-bit word to represent each character. The low byte is an extended ASCII code for the character. (All 255 codes are used, the extra ones are for funny characters like smiley faces and hearts.) The high byte is used for *attributes*. This *attribute byte* is pictured in Figure A.2. The high bit of the attribute byte controls character blinking (it's blinking if the bit is set). The next three bits determine the background color. The next bit controls the intensity of the foreground color—of the character itself. If it's set, the character is displayed in high intensity. The bottom three bits determine the foreground color (the color of the character itself). The only difference between the color and monochrome adapters is the way that the colors are interpreted. The monochrome adapter recognizes only black and white; the code for blue, when used as a foreground color, causes the MGA to print the character underlined.

Note that the foregoing applies only to the CGA and MGA, other adapters do things slightly differently, and the situation is actually more complex for the CGA. If you have a CGA and want more details, video interfacing is covered in depth in Chapter 4 of Peter Norton's book: *The Peter Norton Programmer's Guide to the IBM PC* (Bellevue, Wash.: Microsoft Press, 1985).

The following functions use direct-video access to address the screen:

int dv_init(void)

Initialize the direct-video functions. At present this routine just checks to see if a color card is installed and changes the internal base address of the video memory.

Memory-mapped screens, MGA, CGA.

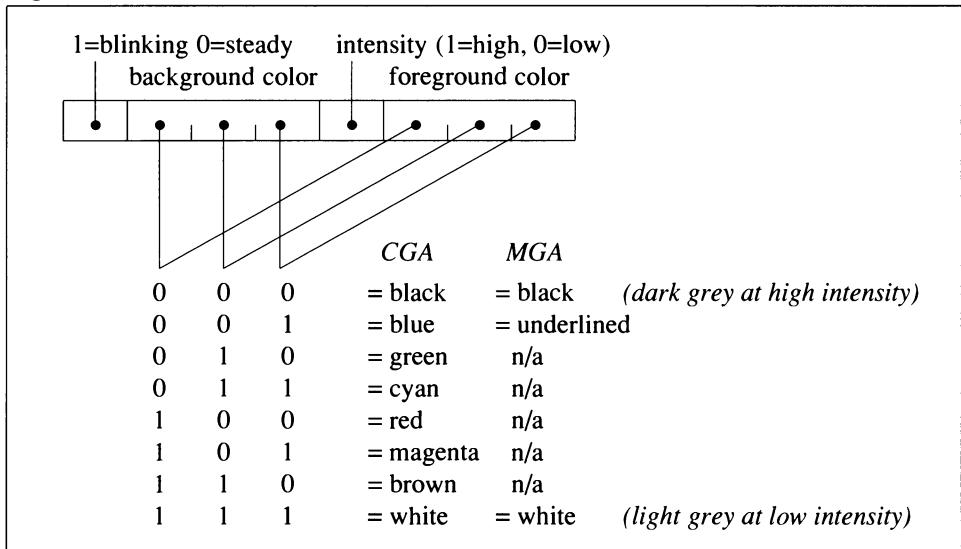
Video display pages.

Characters and attributes.

Blinking, intensity, color.

Direct-video functions.

Initialize direct-video functions.

Figure A.2. Attribute Bits

(The monochrome and color cards map video memory to different base addresses.) It returns zero if an 80-column text mode is not available (the direct-video functions will not work in this case), or 1 if everything's okay.

Clear region of screen.

```
void dv_clr_region(int left, int right, int top,  
                    int bottom, int attrib)
```

Clear the region of screen defined by a box with the upper-left corner at (left,top) and the lower-right corner at (right,bottom). The region is cleared by filling the area with space characters having the indicated attribute. Symbolic defines for the attributes are in <tools/termlib.h>, discussed below. The cursor position is not modified.

Clear screen.

```
void dv_clrs(attrib)
```

Clear the entire screen by filling it with space characters having the indicated attribute. The cursor position is not modified.

Direct-video printf().

```
void dv_printf(int attrib, char *fmt, ...)
```

A printf() that uses direct-video writes. Output characters all have the indicated attribute. The prnt() function described earlier does the actual printing.

Display character and attribute.

```
void dv_putc(int c, int attrib)
```

Write a single character to the screen at the current cursor position, and with the indicated attribute. The following characters are special:

\0	ignored
\f	clear screen and home cursor
\n	send cursor to left edge of next line
\r	send cursor to left edge of current line
\b	back up cursor one character (non-destructive backspace)

The screen scrolls up if you go past the bottom line. Characters that go beyond the end of the current line wrap around to the next line.

void dv_ctoyx(int y, int x)	Change cursor position.
Position the cursor at the indicated row (y) and column (x). The top-left corner of the screen is (0,0), the bottom-right corner is (24,79).	
void dv_getyx(int *rowp, int *colp)	Get cursor position.
Modify *rowp and *colp to hold the current cursor position. The top-left corner of the screen is (0,0), the bottom-right corner is (24,79).	
dv_incha(void)	Input character from screen.
Return the character and attribute at current cursor position. The character is in the low byte of the returned value and the attribute is in the high byte.	
void dv_outcha(int c)	Display character and attribute without moving cursor.
Write character and attribute to screen without moving the cursor. The character is in the low byte of c and the attribute is in the high byte. Note that the special characters supported by dv_putc() are not supported here. They will print as funny-looking IBM graphics characters.	
void dv_replace(int c)	Display character without moving cursor.
Like dv_outcha(), this function writes the character to the screen without moving the cursor, but the attribute is not modified by dv_replace().	
void dv_putchar(int c)	Display normal character.
Write a character to the screen with normal attributes (black background, white foreground, not blinking, normal-intensity, not underlined). This function uses dv_putc() for its output, so it supports the same special characters, scrolling, line wrap, and so forth.	
dv_puts(char *str, int move)	Display string.
Write a string to screen. If move is true, the cursor is positioned at the end of the string, otherwise the cursor is not moved. Normal attributes, as described in earlier dv_putchar(), are used. This function uses dv_putc() for its output.	
void dv_putsa(char *str, int attrib)	Display string with attributes.
Write string to screen, giving characters the indicated attributes. The cursor is positioned just after the rightmost character. This function uses dv_putc() for its output.	
#include <tools/termlib.h>	
SBUF *dv_save (int left, int right, int top, int bottom)	Save and restore region of screen. Discard save buffer.
SBUF *dv_restore (SBUF *sbuf)	
SBUF *dv_freesbuf(SBUF *sbuf)	
These three functions work together to save and restore the characters in a specified region of the screen. dv_save saves all characters in the box with corners at (top, left) and (bottom, right). It returns a pointer that can be passed to a subsequent dv_restore() call to restore that region to its original condition. dv_save terminates the program with an error message if it can't get memory. The cursor is not modified by either function. The SBUF structure, used for this purpose, is defined in <tools/termlib.h>. Note that the memory used for this structure is not freed, you must do that yourself with a df_freesbuf(sbuf) call after the dv_restore() call. For convenience,	

`dv_restore()` returns its own argument, so you can say:

```
    dv_freesbuf( dv_restore( sbuf ) );
```

if you like.

Scroll region arbitrary direction.

```
void dv_scroll_line(int left, int right, int top, int bottom,  
                    int dir, int attrib)
```

Scroll the indicated region of the screen by one line or column. The dir argument must be one of the characters: 'u', 'd', 'l', or 'r' for up, down, left, or right. The cursor is not moved. The opened line or column is filled with space characters having the indicated attribute.

Scroll region up or down.

```
void dv_scroll(int left, int right, int top, int bottom, int amt,  
                int attrib)
```

Scroll the indicated region of the screen up or down by the indicated number of lines (amt), filling the new lines with space characters having the indicated attribute. Negative amounts scroll down, positive amounts scroll up.

Video-BIOS functions that mimic direct-video functions.

There are two sets of I/O routines that use the video-BIOS rather than direct-video reads and writes. The routines in Table A.2 behave exactly like the `dv_` routines that were discussed earlier. For example, `vb_getyx()` works exactly like `dv_getyx()`, but it uses the video BIOS. Names that are in all caps are implemented as macros.

Table A.2. Video-BIOS Macros and Functions That Work Like Direct-Video Functions

Function or macro. [†]			Has side effects?
void	VB_CLRS	(attrib)	no
void	VB_CLR_REGION	(left, right, top, bottom, attrib)	yes
void	VB_CTOYX	(y, x)	no
void	vb_freesbuf	(sbuf)	—
void	vb_getyx	(yp, xp)	—
int	VB_INCHA	()	no
void	VB_OUTCHA	(c)	yes
void	VB_PUTCHAR	(c)	no
void	vb_putc	(c, attrib)	—
void	vb_puts	(str, move_cur)	—
void	VB_REPLACE	(c)	yes
SBUF	*vb_restore	(sbuf)	—
SBUF	*vb_save	(left, right, top, bottom)	—
void	VB_SCROLL	(left, right, top, bottom, amt, attrib)	yes

[†] Names in all caps are macros.

Other video-BIOS functions.

The second set of BIOS routines gives you access to features that are not easy to implement directly. Of these, the cursor-movement functions are used by the direct-video routines to move the physical cursor. As with the earlier group, names in all caps represent macros, but none of these macros have side effects.

Create block cursor.

```
void VB_BLOCKCUR(void)
```

Make the cursor a block cursor rather than the normal underline.

Create underline cursor.

```
void VB_NORMALCUR(void)
```

Make the cursor a normal, underline, cursor.

void VB_CURSIZE(**int** top, **int** bottom)

Change cursor size.

Change the cursor size by making it extend from the indicated top scan line to the bottom one. The line numbers refer to the box in which the character is drawn. A character on the CGA is 8 scan lines (dots) high, and the line numbers go from 0 to 7. They go from 0 to 12 on the MGA. A normal, underline cursor can be created with VB_CURSIZE(6, 7) on the CGA, and VB_CURSIZE(11, 12) on the MGA. VB_CURSIZE(0, 7) creates a block cursor on the CGA, filling the entire area occupied by a character. VB_CURSIZE(0, 1) puts a line over the character rather than under it. VB_CURSIZE(13, 13) makes the cursor disappear entirely. If the top line is larger than the bottom, you'll get a two-part cursor, so VB_CURSIZE(11, 1) creates a cursor with lines both above and below the character on the MGA.

int vb_getchar(**void**)

Get character from keyboard.

Get a character directly from the keyboard. The typed character is returned in the low byte of the returned integer, the high byte holds the auxiliary byte that marks ALT keys and such. See the *IBM Technical Reference* for more information. You can use this function to read key codes that can't be accessed by some compilers' standard I/O system. Similarly, vb_getchar() gets characters from the keyboard, even if input is redirected, though it's more portable to open the console directly for this purpose.

int VB_GETCUR(**void**)

Get cursor position.

Return the current cursor position. The top byte of the return value holds the row, the bottom byte the column.

int VB_GETPAGE(**void**)

Get video page number.

Return a unique number identifying the currently active video-display page.

int vb_iscolor(**void**)

Determine display adapter.

Returns one of the following values:

1	CGA is installed and it's in an 80-column text mode
0	MGA is installed
-1	CGA is installed and it's not in an 80-column text mode

The return value is controlled by the current video mode as follows:

Mode	Return value
2 or 3	1
7	0
anything else	-1

void VB_SETCUR(**posn**)

Move physical cursor.

Modify current cursor position. The top byte of posn holds the row (y), the bottom byte, the column (x). The top-left corner of the screen is (0,0).

A.9.2 Video I/O—Implementation

This section presents the code that implements the foregoing functions. Most of the actual code is pretty boring, and it's presented with no additional comment in the text. A high-level discussion of the data structures and techniques used is discussed however.

The `int86()` function.

Attribute definitions,
`SBUF`, `WORD`, `IMALLOC`,
`IFREE`: `<tools/termlib.h>`

I'm assuming, in this section, that you're familiar to some extent with the IBM ROM-BIOS interface. If you need this information, read Peter Norton's book: *The Peter Norton Programmer's Guide to the IBM PC* (Bellevue, Wash.: Microsoft Press, 1985), which covers all this material. I'm also assuming that your compiler has a mechanism for generating a software interrupt—most do. I'm using the Microsoft `int86()` because it's portable—newer versions of the compiler have more efficient BIOS-access functions in the library, and you may want to modify the code to use these functions.

The `<tools/termlib.h>` file is used both by application programs and by the I/O functions themselves. It's shown in Listing A.42. The macros on lines five to 12 are the codes for the basic colors described earlier. The `FGND()` and `BGND()` macros on lines 14 and 15 put the color code into the foreground or background position within the attribute byte. (`FGND` is just for documentation, it doesn't do anything. `BGND` shifts the color four bits to the left.) The `NORMAL`, `UNDERLINED`, and `REVERSE` definitions on lines 17 to 19 of Listing A.42 define all legal color combinations for the monochrome card: a normal character, an underlined character, and a reverse-video character. Finally, `BLINKING` and `BOLD` (on lines 21 and 22 of Listing A.42) can be ORed with the colors to make the character blink or be displayed at high intensity.

Listing A.42. *termlib.h*—Video-I/O Definitions

```

1  /* Various definitions for the termlib. Note that if your program includes both
2   * termlib.h and vbios.h, termlib.h must be included FIRST.
3   */
4
5 #define BLACK          0x00      /* Color Card.      */
6 #define BLUE           0x01
7 #define GREEN          0x02
8 #define CYAN           0x03
9 #define RED            0x04
10 #define MAGENTA        0x05
11 #define BROWN          0x06
12 #define WHITE          0x07
13
14 #define FGND(color)    (color)
15 #define BGND(color)    ((color) <<4)
16
17 #define NORMAL         (FGND(WHITE) | BGND(BLACK))  /* Monochrome card */
18 #define UNDERLINED     (FGND(BLUE)  | BGND(BLACK))
19 #define REVERSE        (FGND(BLACK) | BGND(WHITE))
20
21 #define BLINKING       0x80      /* May be ORed with the above */
22 #define BOLD            0x08      /* and with each other */
23
24 /-----
25 * If USE_FAR_HEAP is true then use the far heap to save screen images in the
26 * small model. You must recompile the termlib if you change this #define.
27 */
28
29 typedef unsigned int WORD;
30
31 #if( USE_FAR_HEAP )
32     typedef WORD far *IMAGEP;
33     #define IMALLOC _fmalloc
34     #define IFREE _ffree
35 #else
36     typedef WORD     *IMAGEP;
37     #define IMALLOC malloc
38     #define IFREE   free

```



Listing A.42. continued...

```

39 #endif
40
41 typedef struct SBUF /* used by vb_save, vb_restore, dv_save, and dv_restore */
42 {
43     unsigned int top, bottom, left, right;
44     IMAGEP image;
45 } SBUF;
46
47 /-----
48 *      Prototypes for the video-BIOS access routines.
49 */
50
51 extern int vb_iscolor      ( void );
52 extern void vb_getyx      ( int *yp, int *xp );
53 extern void vb_putc       ( int c, int attrib );
54 extern void vb_puts       ( char *str, int move_cur );
55 extern int vb_getchar    ( void );
56 extern SBUF *vb_save     ( int l, int r, int t, int b );
57 extern SBUF *vb_restore   ( SBUF *sbuf );
58 extern void vb_freesbuf  ( SBUF *sbuf );
59
60 /-----
61 *      Prototypes for the equivalent direct video functions.
62 */
63
64 extern int dv_init        ( void );
65 extern void dv_scroll_line ( int x_left, int x_right, int y_top,
66                                int y_bottom, int dir, int attrib );
67 extern void dv_scroll     ( int x_left, int x_right, int y_top,
68                                int y_bottom, int amt, int attrib );
69 extern void dv_clrs       ( int attrib );
70 extern void dv_clr_region ( int l, int r, int t, int b, int attrib );
71 extern void dv_ctoyx      ( int y, int x );
72 extern void dv_getyx      ( int *rowp, int *colp );
73 extern void dv_putc       ( int c, int attrib );
74 extern void dv_putchar    ( int c );
75 extern int dv_puts       ( char *str, int move_cur );
76 extern void dv_putsa     ( char *str, int attrib );
77 extern int dv_incha     ( void );
78 extern void dv_outcha    ( int c );
79 extern void dv_replace    ( int c );
80 extern void dv_printf    ( int attribute, char *fmt, ... );
81 extern SBUF *dv_save     ( int l, int r, int t, int b );
82 extern SBUF *dv_restore   ( SBUF *sbuf );
83 extern void dv_freesbuf  ( SBUF *sbuf );

```

The WORD **typedef** on line 29 of Listing A.42 is an attempt to get a portable definition for a 16-bit unsigned quantity. It's used later to access a single 16-bit, character/attribute pair from the video page.

Portable 16-bit data type, WORD.

The definitions on lines 31 to 39 of Listing A.42 are compiler dependent. They control from where, in real memory, the buffers used for an SBUF will be allocated. In an average window application, where you need to save the area under a new window before creating it, a considerable amount of memory can be used to store these old images ($25 \times 80 \times 16 = 4,000$ bytes for the whole screen), and that much memory may not be available in an 8086 small-model program. You don't want to convert the entire program to the medium or compact model because this conversion will slow down all the

Screen-save buffer, SBUF.

The ‘far heap,’
`USE_FAR_HEAP,`
`_fmalloc(),`
`_ffree().`

The `far` keyword.

Video-BIOS definitions,
`vbios.h`.

pointer accesses. It’s possible, however, for a small-model program to allocate and use a region of memory outside of the normal 64K data area (called the *far heap*). The Microsoft `_fmalloc()` function gets memory from the far heap, and returns a 32-bit far pointer. The `_ffree()` function puts memory back into the far heap. The far heap is used for window images if `USE_FAR_HEAP` is defined above line 31 of Listing A.42 (or with a `-DUSE_FAR_HEAP` command-line switch). The `far` keyword on line 32 tells the Microsoft compiler to use a 32-bit pointer that holds both the segment and offset components of the address when `_fmalloc()` is used to allocate memory.

The next `.h` file of interest is `vbios.h`, which is used locally by the video-BIOS functions. It’s in Listing A.43. Two interrupts are of interest here: the video interrupt, whose number is defined on line five, and the keyboard interrupt defined on the next line. The definitions on lines seven to 15 let you select a specific function—each interrupt can do several things, and the function-number determines which of these functions are performed. All of the foregoing is described in great depth in Peter Norton’s book if you need more background information.

Listing A.43. `vbios.h`— Video-BIOS Service Definitions

```

1  #ifndef NORMAL
2  #include <tools/termlib.h>
3  #endif
4
5  #define VIDEO_INT      0x10    /* Video interrupt          */
6  #define KB_INT         0x16    /* Keyboard interrupt       */
7  #define CUR_SIZE        0x1     /* Set cursor size          */
8  #define SET_POSN        0x2     /* Modify cursor posn      */
9  #define READ_POSN       0x3     /* Read current cursor posn */
10 #define SCROLL_UP        0x6     /* scroll region of screen up */
11 #define SCROLL_DOWN      0x7     /* " down                  */
12 #define READ_CHAR        0x8     /* Read character from screen */
13 #define WRITE           0x9     /* Write character          */
14 #define WRITE_TTY        0xe     /* Write char & move cursor */
15 #define GET_VMODE        0xf     /* Get video mode & disp pg */

```

Accessing the video
 BIOS, `_Vbios()`.

Most video-BIOS functions implemented as
 macros.

Direct-video definitions,
`video.h`. Video-memory
 base addresses, `MON-`
`BASE`, `COLBASE`.

Screen dimensions, `NUM-`
`ROWS`, `NUMCOLS`.

Accessing characters
 and attributes, `CHARAC-`
`TER`. Representing entire
 screen, `DISPLAY`.

The video BIOS is accessed by `_Vbios()` in Listing A.44. The Microsoft `int86()` function, a prototype for which is in `<dos.h>`, communicates with the BIOS using the normal interrupt mechanism. It is passed an interrupt number and pointers to two REGS structures that represent the 8086 register set. The first of these holds the desired contents of the registers before the interrupt is executed, the second holds the values of the registers after the interrupt returns. Most of the video-BIOS functions are actually macros that evaluate to `_Vbios()` calls. These are also defined in `<tools/vbios.h>` and are shown in Listing A.45.

The third `.h` file, `video.h` in Listing A.46, is used locally by the direct-video functions. `MONBASE` and `COLBASE` (defined on lines one and two of Listing A.46) are the base addresses of page 0 of the monochrome and color adapters. The addresses are in the canonical segment/offset form used by most 8086 compilers, the high 16 bits are the segment and the low 16 bits are the offset. The dimensions of the screen (in characters) are controlled by `NUMROWS` and `NUMCOLS` on the lines three and four of Listing A.46.

The `CHARACTER` structure defined on lines six to 11 of Listing A.46 describes a single character/attribute pair. Using a structure to access the high byte is a better strategy than a shift because it’s more efficient in most compilers. A `DISPLAY`, defined on the next line, is a 25×80 array of these character/attribute pairs.

Listing A.44. *_vbios.c*— Video-BIOS Interface Function

```

1 #include <dos.h>
2 #include <tools/vbios.h>
3 #include "video.h"
4
5 /* This file contains a workhorse function used by the other video I/O */
6 /* functions to talk to the BIOS. It executes the video interrupt. */
7
8 int _Vbios( service, al, bx, cx, dx, return_this )
9 int service;           /* Service code, put into ah. */
10 int al, bx, cx, dx;    /* Other input registers. */
11 char *return_this;     /* Register to return, "ax", "bh", "dx" only. */
12 {
13     union REGS regs;
14
15     regs.h.ah = service;
16     regs.h.al = al;
17     regs.x.bx = bx;
18     regs.x.cx = cx;
19     regs.x.dx = dx;
20     int86( VIDEO_INT, &regs, &regs );
21     return (*return_this == 'a') ? regs.x.ax :
22         (*return_this == 'b') ? regs.h.bh : regs.x.dx ;
23 }
```

The `dv_Screen` pointer declared on line 16 of Listing A.46 points at the base address of the video memory for the current adapter (it's initialized for the monochrome card but it can be changed at run time). The `SCREEN` macro on line 21 of Listing A.46 uses `dv_Screen` to randomly access a specific character or attribute on the screen. For example, the following puts a high-intensity X in the lower-right corner of the screen:

```

#include <termlib.h>
#include <video.h>
SCREEN[ 25 ][ 79 ].character = 'X' ;
SCREEN[ 25 ][ 79 ].attribute = NORMAL | BOLD ;
```

The elaborate casts are needed because you can't cast a pointer into an array, but you can cast it to an array pointer. Since `dv_Screen` is a pointer to an entire two-dimensional array as compared to a pointer to the first element, the star in `SCREEN` gets you the array itself. Looked at another way, a pointer to an entire array is special in two ways. First, if you increment it, you'll skip past the entire array, not just one element. Second, the pointer is treated internally as if it were a pointer to a pointer. That is, the array pointer is treated internally as if it points at an array of pointers to rows, similar to `argv`. An extra star is needed to select the correct row. This star doesn't change the value of the pointer, only its type—`dv_Screen` and `*dv_Screen` both evaluate to the same number, but the types are different. `dv_Screen` (without the star) is of type *pointer to two-dimensional array*, the first element of which is one-dimensional array of `CHARACTERS`. `*dv_Screen` (with the star) is of type *pointer to one-dimensional array (pointer to row)*, the first element of which is a single `CHARACTER`. All this is discussed in greater depth in Chapter 6.

The definitions on lines 23 to 25 of Listing A.46 work just like the foregoing, except that the screen is treated as an array of 16-bit numbers rather than as an array of `CHARACTERS`. The `dv_` functions are in Listings A.47 to A.67.

Accessing video memory directly, `dv_Screen`, `SCREEN`

Listing A.45. *vbios.h*— Macros That Implement Video-BIOS Functions

```

16 /* These video-BIOS functions are implemented as macros.
17 *
18 * VB_INCHA      Returns the character and attribute ORed together. Character in
19 *                the low byte and attribute in the high byte.
20 * VB_GETPAGE    Return the currently active display page number
21 * VB_GETCUR     Get current cursor position. The top byte of the return value
22 *                holds the row, the bottom by the column. Pagenum is the video
23 *                page number. Note that VB_GETPAGE() will mess up the fields in
24 *                the Regs structure so it must be called first.
25 * VB_CURSIZE   Change the cursor shape to go from the top to the bottom scan
26 *                line.
27 * VB_OUTCHA   Write a character and attribute without moving the cursor. The
28 *                attribute is in c's high byte, the character is the low byte.
29 * VB_REPLACE    Same as VB_OUTCHA but uses the existing attribute byte.
30 * VB_SETCUR    Modify current cursor position. The top byte of "posn" value
31 *                holds the row (y), the bottom byte, the column (x). The top-left
32 *                corner of the screen is (0,0). Pagenum is the video-display-page
33 *                number.
34 * VB_CTOYX(y,x) Like VB_SETCUR but y and x coordinates are used.
35 * VB_SCROLL     Scroll the indicated region on the screen. If amt is <0,
36 *                scroll down; otherwise, scroll up.
37 * VB_CLRS       Clear the entire screen
38 * VB_CLR_REGION Clear a region of the screen
39 * VB_BLOCKCUR  Change to a block cursor.
40 * VB_NORMALCUR Change to an underline cursor.
41 * VB_PUTCHAR   Like vb_putc, but uses white on black for the attribute.
42 */
43
44 #define VB_GETPAGE()      _Vbios( GET_VMODE, 0, 0, 0, 0,                                "bh")
45 #define VB_INCHA()       _Vbios( READ_CHAR, 0, VB_GETPAGE(), 0, 0,                            "ax")
46 #define VB_GETCUR()      _Vbios( READ_POSN, 0, VB_GETPAGE(), 0, 0,                            "dx")
47 #define VB_CURSIZE(t,b)  _Vbios( CUR_SIZE, 0,0,((t)<<8)|(b),0,                           "ax")
48 #define VB_OUTCHA(c)    _Vbios( WRITE,      (c)&0xff,((c)>>8)&0xff, 1, 0, "ax")
49 #define VB_REPLACE(c)   VB_OUTCHA( (c & 0xff) | (VB_INCHA() & ~0xff) )
50 #define VB_SETCUR(posn) _Vbios( SET_POSN, 0, VB_GETPAGE() << 8, 0,(posn), "ax")
51 #define VB_CTOYX(y,x)   VB_SETCUR( ((y) << 8) | ((x) & 0xff) )
52 #define VB_SCROLL(xl, xr, yt, yb, amt, attr) _Vbios( \
53                                     ((amt) < 0) ? SCROLL_DOWN : SCROLL_UP, \
54                                     abs(amt), (attr) << 8, ((yt) << 8) | (xl), \
55                                     ((yb) << 8) | (xr), "ax"\ \
56                                     )
57
58 #define VB_CLRS(at)      VB_SCROLL( 0, 79, 0, 24, 25,                                (at))
59 #define VB_CLR_REGION(l,r,t,b,at) VB_SCROLL( (l),(r),(t),(b),((b)-(t))+1,(at))
60
61 #define VB_BLOCKCUR()   VB_CURSIZE( 0, vb_iscolor() ? 7 : 12 )
62 #define VB_NORMALCUR() ( vb_iscolor() ? VB_CURSIZE(6,7) : VB_CURSIZE(11,12) )
63 #define VB_PUTCHAR(c)  vb_putc( (c), NORMAL )

```

Listing A.46. *video.h*—Definitions for Direct-Video Functions

```

1 #define MONBASE          (DISPLAY far *) 0xb0000000
2 #define COLBASE          (DISPLAY far *) 0xb8000000
3 #define NUMROWS           25
4 #define NUMCOLS           80
5
6 typedef struct
7 {
8     unsigned char letter;
9     unsigned char attribute;
10 } CHARACTER;
11
12 typedef CHARACTER      DISPLAY[ NUMROWS ][ NUMCOLS ];
13
14 #ifdef ALLOC
15     DISPLAY far *dv_Screen = (DISPLAY far *) MONBASE ;
16 #else
17     extern DISPLAY far *dv_Screen;
18 #endif
19
20 #define SCREEN          (* dv_Screen)
21
22 typedef short           CHAR_ATTRIB;
23 typedef CHAR_ATTRIB    VDISPLAY[ NUMROWS ][ NUMCOLS ];
24 #define VSCREEN          (* (VDISPLAY far *)dv_Screen )

```

Listing A.47. *dv_clr_r.c*—Clear Region of Screen (Direct Video)

```

1 #include "video.h"
2
3 void     dv_clr_region( l, r, t, b, attrib )
4 {
5     int      ysize, xsize, x, y ;
6
7     xsize = (r - l) + 1;                      /* horizontal size of region */
8     ysize = (b - t) + 1;                      /* vertical size of region */
9
10    for( y = ysize; --y >= 0 ; )
11        for( x = xsize; --x >= 0 ; )
12        {
13            SCREEN[ y + t ][ x + l ].letter    = ' ';
14            SCREEN[ y + t ][ x + l ].attribute = attrib ;
15        }
16 }

```

Listing A.48. *dv_clrs.c*—Clear Entire Screen (Direct Video)

```

1 #include "video.h"
2
3 void     dv_clrs( attrib )
4 {
5     /* Clear the screen. The cursor is not moved. */
6

```



Listing A.48. continued...

```

7   CHARACTER far      *p = (CHARACTER far *)( dv_Screen );
8   register int        i ;
9
10  for( i = NUMROWS * NUMCOLS; --i >= 0 ; )
11  {
12      (p )->letter    = ' ' ;
13      (p++)->attribute = attrib ;
14  }
15 }
```

Listing A.49. dv_frees.c— Free an SBUF (Direct Video)

```

1 #include <stdlib.h>
2 #include <tools/termlib.h>
3 #include "video.h"
4                                         /* Free an SBUF as is allocated by vb_save(). */
5 void     dv_freesbuf( p )
6 SBUF     *p;
7 {
8     IFREE( p->image );
9     free ( p );
10 }
```

Listing A.50. dv_init.c— Initialize Direct-Video Functions

```

1 #include "video.h"
2
3 int     dv_init()
4 {
5     int i;
6     if( (i = vb_iscolor()) >= 0 )
7         dv_Screen = i ? COLBASE : MONBASE ;
8     return( i != -1 );
9 }
```

Listing A.51. dv_printf.c— A Direct-Video printf()

```

1 #include <stdio.h>
2 #include <stdarg.h>
3 #include <tools/debug.h>          /* For VA_LIST definition */
4 #include "video.h"
5
6 void     dv_printf( attrib, fmt, VA_LIST )
7 int     attrib;
8 char    *fmt;
9 {
10    /* Direct-video printf, characters will have the indicated attributes. */
11    /* prnt() is in curses.lib, which must be linked to your program if you */
12    /* use the current function. */
13 }
```



Listing A.51. continued...

```

14     extern dv_putc();
15     va_list args;
16
17     va_start( args, fmt );
18     prnt    ( dv_putc, attrib, fmt, args );
19     va_end  ( args );
20 }
```

Listing A.52. dv_putc.c—Write Character to Screen with Attribute (Direct Video)

```

1  #include "video.h"
2  #include <tools/vbios.h>
3
4  static int      Row = 0;           /* Cursor Row      */
5  static int      Col = 0;           /* Cursor Column   */
6
7  /*-----*/
8  /* Move cursor to (Row, Col) */
9
10 #define fix_cur() _Vbios( SET_POSN, 0, 0, 0, (Row << 8) | Col , "ax" )
11
12 /*-----*/
13
14 void      dv_putc( c, attrib )
15 {
16     /* Write a single character to the screen with the indicated attribute.
17      * The following are special:
18      *
19      * \0      ignored
20      * \f      clear screen and home cursor
21      * \n      to left edge of next line
22      * \r      to left edge of current line
23      * \b      one character to left (non-destructive)
24      *
25      * The screen will scroll up if you go past the bottom line. Characters
26      * that go beyond the end of the current line wrap around to the next line.
27      */
28
29     switch( c )
30     {
31         case 0:      break;          /* Ignore ASCII NULL's */
32
33         case '\f':  dv_clrs( attrib );
34             Row = Col = 0;
35             break;
36
37         case '\n':  if( ++Row >= NUMROWS )
38         {
39             dv_scroll_line(0,79,0,24, 'u', NORMAL );
40             Row = NUMROWS-1 ;
41         }
42             /* Fall through to '\r' */
43
44         case '\r':  Col = 0;
45         break;
46 }
```



Listing A.52. continued...

```

46     case '\b': if( --Col < 0 )
47         Col = 0;
48         break;
49
50     default : SCREEN[ Row ][ Col ].letter      = c      ;
51             SCREEN[ Row ][ Col ].attribute = attrib ;
52             if( ++Col >= NUMCOLS )
53             {
54                 Col = 0;
55                 if( ++Row >= NUMROWS )
56                 {
57                     dv_scroll_line(0,79,0,24, 'u', NORMAL );
58                     Row = NUMROWS-1 ;
59                 }
60             }
61             break;
62     }
63     fix_cur();
64 }
65
66 /*-----*/
67
68 void dv_ctoyx( y, x )           /* Position the cursor at the indicated row and column */
69 {
70     Row = y;
71     Col = x;
72     fix_cur();
73 }
74
75 /*-----*/
76
77 void    dv_getyx( rowp, colp )
78 int     *rowp, *colp;
79 {                                /* Modify *rowp and *colp to hold the cursor position. */
80     *rowp = Row;
81     *colp = Col;
82 }
83
84 /*-----*/
85
86 int     dv_incha()            /* Get character & attribute from screen. */
87 {
88     return (int) VSCREEN[ Row ][ Col ];
89 }
90
91 void    dv_outcha(c)          /* Write char. & attrib. w/o moving cursor. */
92 {
93     VSCREEN[ Row ][ Col ] = c ;
94 }
95
96 void    dv_replace(c)         /* Write char. only w/o moving cursor */
97 {
98     SCREEN[ Row ][ Col ].letter = c ;
99 }
```

Listing A.53. dv_putchar.c— Write Character to Screen, Normal Attribute (Direct Video)

```
1 #include "video.h"
2 #include <tools/termlib.h>
3
4 void     dv_putchar( c )
5 {
6     dv_putc( c & 0xff, NORMAL );
7 }
```

Listing A.54. dv_puts.c— Write String to Screen, Normal Attribute (Direct Video)

```
1 #include "video.h"
2 #include <tools/termlib.h>
3
4 dv_puts( str, move )
5 char      *str;
6 {
7     /* Write string to screen, moving cursor to end of string only if move is
8      * true. Use normal attributes.
9      */
10
11    int orow, ocol;
12
13    dv_getyx( &orow, &ocol );
14
15    while( *str )
16        dv_putc( *str++, NORMAL );
17
18    if( !move )
19        dv_ctoyx( orow, ocol );
20 }
```

Listing A.55. dv_putsa.c— Write String to Screen, with Attribute (Direct Video)

```
1 #include "video.h"
2
3 void     dv_putsa( str, attrib )
4 register char   *str;
5 register int    attrib;
6 {
7     /* Write string to screen, giving characters the indicated attributes. */
8
9     while( *str )
10        dv_putc( *str++, attrib );
11 }
```

Listing A.56. dv_resto.c—Restore Saved Region (Direct Video)

```

1 #include <stdio.h>
2 #include <tools/termlib.h>
3 #include "video.h"
4
5 SBUF *dv_restore( sbuf )
6 SBUF *sbuf;
7 {
8     /* Restore a region saved with a previous dv_save() call. The cursor is
9      * not modified. Note that the memory used by sbuf is not freed, you must
10     * do that yourself with a dv_freesbuf(sbuf) call.
11     */
12
13     int ysize, xsize, x, y ;
14     IMAGEP p;
15
16     xsize = ( sbuf->right - sbuf->left ) + 1 ;
17     ysize = ( sbuf->bottom - sbuf->top ) + 1 ;
18     p = sbuf->image;
19
20     for( y = 0; y < ysize ; ++y )
21         for( x = 0; x < xsize ; ++x )
22             VSCREEN[ y + sbuf->top ][ x + sbuf->left ] = *p++;
23
24     return sbuf;
25 }
```

Listing A.57. dv_save.c—Save Region (Direct Video)

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include "video.h"
4 #include <tools/termlib.h>
5
6 SBUF *dv_save( l, r, t, b )
7 {
8     /* Save all characters and attributes in indicated region. Return a
9      * pointer to a save buffer. The cursor is not modified. Note that the
10     * save buffer can be allocated from the far heap, but the SBUF itself is
11     * not. See also, dv_restore() and dv_freesbuf(); */
12
13
14     int ysize, xsize, x, y ;
15     IMAGEP p;
16     SBUF *sbuf;
17
18     xsize = (r - l) + 1;
19     ysize = (b - t) + 1;
20
21     if( !(sbuf = (SBUF *) malloc(sizeof(SBUF)) ) )
22     {
23         fprintf(stderr, "Internal error (dv_save): No memory for SBUF.");
24         exit( 1 );
25     }
26     if( !(p = (IMAGEP) IMALLOC(xsize * ysize * sizeof(WORD))) )
27     {
28         fprintf(stderr, "Internal error (dv_save): No memory for image");
29         exit( 2 );

```

Listing A.57. continued...

```

30     }
31
32     sbuf->left    = l;
33     sbuf->right   = r;
34     sbuf->top     = t;
35     sbuf->bottom  = b;
36     sbuf->image   = p;
37
38     for( y = 0; y < ysize ; ++y )
39         for( x = 0; x < xsize ; ++x )
40             *p++ = VSCREEN[ y + t ][ x + l ];
41
42     return sbuf;
43 }
```

Listing A.58. dv_scree.c— Direct-Video Variable Allocation

```

1 #define ALLOC
2 #include "video.h"
3
4 /* This file has no code in it. It just allocates space for the variables
5  * defined in video.h
6 */

```

Listing A.59. dv_scrol.c— Scroll Region of Screen (Direct Video)

```

1 #include "video.h"
2
3 static void cpy_row( dest_row, src_row, left_col, right_col )
4 {
5     /* Copy all characters between left_col and right_col (inclusive)
6      * from src_row to the equivalent position in dest_row.
7      */
8
9     CHARACTER far *s ;
10    CHARACTER far *d ;
11
12    d = & SCREEN[ dest_row ][ left_col ];
13    s = & SCREEN[ src_row ][ left_col ];
14
15    while( left_col++ <= right_col )
16        *d++ = *s++;
17 }
18
19 /*-----*/
20
21 static void cpy_col( dest_col, src_col, top_row, bot_row )
22 {
23     /* Copy all characters between top_row and bot_row (inclusive)
24      * from src_col to the equivalent position in dest_col.
25      */
26
27     CHARACTER far *s = & SCREEN[ top_row ][ src_col ];
28     CHARACTER far *d = & SCREEN[ top_row ][ dest_col ];
29 }
```



Listing A.59. continued...

```

30     while( top_row++ <= bot_row )
31     {
32         *d = *s;
33         d += NUMCOLS;
34         s += NUMCOLS;
35     }
36 }
37
38 /*-----*/
39
40 static void clr_row( row, attrib, left_col, right_col )
41 {
42     /* Clear all characters in the indicated row that are between left_col and
43      * right_col (inclusive).
44      */
45
46     CHARACTER far *p = & SCREEN[ row ][ left_col ];
47
48     while( left_col++ <= right_col )
49     {
50         (p )->letter    = ' ';
51         (p++)->attribute = attrib ;
52     }
53 }
54
55 /*-----*/
56
57 static void clr_col( col, attrib, top_row, bot_row )
58 {
59     /* Clear all characters in the indicated column that are between top_row
60      * and bot_row (inclusive).
61      */
62
63     CHARACTER far *p = & SCREEN[ top_row ][ col ];
64
65     while( top_row++ <= bot_row )
66     {
67         p->letter    = ' ';
68         p->attribute = attrib ;
69         p += NUMCOLS ;
70     }
71 }
72
73 /*=====
74      *          Externally accessible functions:
75      *=====
76 */
77
78 void dv_scroll_line ( x_left, x_right, y_top, y_bottom, dir, attrib )
79 {
80     /* Scroll the window located at:
81      *
82      * (y_top, x_left)
83      *           +-----+
84      *           |       |
85      *           |       |
86      *           +-----+
87      *                   (y_bottom, x_right)
88      *
89      * Dir is one of: 'u', 'd', 'l', or 'r' for up, down, left, or right.

```



Listing A.59. continued...

```
90     * The cursor is not moved. The opened line is filled with space characters
91     * having the indicated attribute.
92     */
93
94     int      i;
95     CHARACTER far *p;
96
97     if( dir == 'u' )
98     {
99         for( i = y_top; i < y_bottom ; i++ )
100             cpy_row( i, i+1, x_left, x_right );
101             clr_row( y_bottom, attrib, x_left, x_right );
102     }
103     else if( dir == 'd' )
104     {
105         for( i = y_bottom; --i >= y_top ; )
106             cpy_row( i+1, i, x_left, x_right );
107             clr_row( y_top, attrib, x_left, x_right );
108     }
109     else if( dir == 'l' )
110     {
111         for( i = x_left; i < x_right; i++ )
112             cpy_col( i, i+1, y_top, y_bottom );
113             clr_col( x_right, attrib, y_top, y_bottom );
114     }
115     else /* dir == 'r' */
116     {
117         for( i = x_right; --i >= x_left ; )
118             cpy_col( i+1, i, y_top, y_bottom );
119             clr_col( x_left, attrib, y_top, y_bottom );
120     }
121 }
122
123 /*-----*/
124
125 void dv_scroll( x_left, x_right, y_top, y_bottom, amt, attrib )
126 {
127     /* Scroll the screen up or down by the indicated amount. Negative
128     * amounts scroll down.
129     */
130
131     int dir = 'u';
132
133     if( amt < 0 )
134     {
135         amt = -amt;
136         dir = 'd' ;
137     }
138     while( --amt >= 0 )
139         dv_scroll_line( x_left, x_right, y_top, y_bottom, dir, attrib );
140 }
```

Listing A.60. *vb_frees.c*—Free Save Buffer (Video BIOS)

```

1 #include "video.h"
2 #include <tools/termlib.h>
3 #include <stdlib.h>
4
5 void vb_freesbuf( p )
6 SBUF *p;
7 { /* Free an SBUF as is allocated by vb_save(). */
8     IFREE( p->image );
9     free ( p );
10 }
```

Listing A.61. *vb_getch.c*—Get Character from Keyboard (Keyboard!BIOS)

```

1 #include <dos.h>
2 #include <tools/vbios.h>
3
4 int vb_getchar()
5 {
6     /* Get a character directly from the keyboard */
7
8     union REGS regs;
9     regs.h.ah = 0 ;
10    int86( KB_INT, &regs, &regs );
11    return( (int)regs.x.ax );
12 }
```

Listing A.62. *vb_getyx.c*—Get Cursor Position (Video BIOS)

```

1 #include <tools/vbios.h>
2
3 void vb_getyx( yp, xp )
4 int *yp, *xp;
5 {
6     register int posn;
7
8     posn = VB_GETCUR();
9     *xp = posn & 0xff ;
10    *yp = (posn >> 8) & 0xff ;
11 }
```

Listing A.63. *vb_iscol.c*—Test for Color Card (Video BIOS)

```

1 #include <tools/vbios.h>
2
3 int vb_iscolor() /* Returns true if a color card is active */
4 {
5     int mode = _Vbios( GET_VMODE, 0, 0, 0, 0, "ax" ) & 0xff ;
6     return( (mode == 7) ? 0 :
7            (mode == 2 || mode == 3) ? 1 : -1 );
8 }
```

Listing A.64. *vb_putc.c*—Write Character in TTY Mode (Video BIOS)

```

1 #include <tools/termlib.h>
2 #include <tools/vbios.h>
3
4 void vb_putc( c, attrib )
5 {
6     /* Write a character to the screen in TTY mode. Only normal printing
7      * characters, BS, BEL, CR and LF are recognized. The cursor is automatic-
8      * ally advanced and lines will wrap. The WRITE_TTY BIOS service doesn't
9      * handle attributes correctly, so printing characters have to be output
10     * twice---once by VB_OUTCHA to set the attribute bit, and again using
11     * WRITE_TTY to move the cursor. WRITE_TTY picks up the existing attribute.
12     */
13
14     if( c != '\b' && c != '\007' && c != '\r' && c != '\n' )
15         VB_OUTCHA( (c & 0xff) | (attrib << 8) );
16
17     _Vbios( WRITE_TTY, c, attrib & 0xff, 0, 0, "ax" );
18 }
```

Listing A.65. *vb_puts.c*—Write String in TTY Mode (Video BIOS)

```

1 #include <tools/vbios.h>
2
3 void    vb_puts( str, move_cur )
4 register char  *str;
5 {
6     /* Write a string to the screen in TTY mode. If move_cur is true the cursor
7      * is left at the end of string. If not the cursor will be restored to its
8      * original position (before the write).
9     */
10
11     int posn;
12
13     if( !move_cur )
14         posn = VB_GETCUR();
15
16     while( *str )
17         VB_PUTCHAR( *str++ );
18
19     if( !move_cur )
20         VB_SETCUR( posn );
21 }
```

Listing A.66. *vb_resto.c*—Restore Saved Region (Video BIOS)

```

1 #include "video.h"
2 #include <tools/vbios.h>
3 #include <stdlib.h>
4
5 SBUF    *vb_restore( sbuf )
6 SBUF    *sbuf;
7 {
```



Listing A.66. continued...

```

8  /* Restore a region saved with a previous vb_save() call. The cursor is
9   * not modified. Note that the memory used by sbuf is not freed, you must
10  * do that yourself with a vb_freesbuf(sbuf) call.
11  */
12
13  int      ysize, xsize, x, y ;
14  IMAGEP   p;
15
16  xsize = ( sbuf->right - sbuf->left ) + 1 ;
17  ysize = ( sbuf->bottom - sbuf->top ) + 1 ;
18  p     = sbuf->image;
19
20  for( y = 0; y < ysize ; ++y )
21    for( x = 0; x < xsize ; ++x )
22    {
23      VB_CTOYX( y + sbuf->top, x + sbuf->left );
24      VB_OUTCHA( *p );
25      ++p;           /* VB_OUTCHA has side effects so can't use *p++ */
26    }
27  return sbuf;
28 }
```

Listing A.67. vb_save.c—Save Region (Video BIOS)

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/termlib.h> /* must be included first */
4 #include <tools/vbios.h>
5
6 SBUF *vb_save( l, r, t, b )
7 {
8   /* Save all characters and attributes in indicated region. Return a pointer
9    * to a save buffer. The cursor is not modified. Note that the save buffer
10   * is allocated from the far heap but the sbuf itself is not. See also,
11   * dv_restore() and dv_freesbuf();
12   */
13
14   int      ysize, xsize, x, y ;
15   IMAGEP   p;
16   SBUF     *sbuf;
17
18   xsize = (r - l) + 1;
19   ysize = (b - t) + 1;
20
21   if( !(sbuf = (SBUF *) malloc(sizeof(SBUF)) ) )
22   {
23     fprintf(stderr, "Internal error (vb_save): No memory for SBUF.");
24     exit( 1 );
25   }
26
27   if( !(p = (IMAGEP) IMALLOC(xsize * ysize * sizeof(WORD)))) 
28   {
29     fprintf(stderr, "Internal error (vb_save): No memory for image.");
30     exit( 2 );
31   }
32 }
```



Listing A.67. continued...

```

33     sbuf->left    = l;
34     sbuf->right   = r;
35     sbuf->top     = t;
36     sbuf->bottom  = b;
37     sbuf->image   = p;
38
39     for( y = 0; y < ysize ; ++y )
40         for( x = 0; x < xsize ; ++x )
41         {
42             VB_CTOYX( y + t , x + 1 );
43             *p++ = VB_INCHA();
44         }
45
46     return sbuf;
47 }
```

A.10 Low-level-I/O, Glue Functions

“Glue” functions are intermediate-level functions that allow your program to access low-level functions in a portable way. In the current context, the glue functions let you access either the direct-video or video-BIOS low-level-I/O functions without having to change the function calls in the program itself. If your program uses these glue functions instead of the actual I/O subroutines, you can choose which set of functions to use by linking either direct-video or BIOS libraries to your final program, without having to change the program itself. The glue-to-I/O-function mapping is summarized in Table A.3.

Table A.3. Mapping Glue Routines to I/O Functions

Glue Function†	Mapped to This Function	
	Direct-video mode	BIOS mode
void clr_region(l,r,t,b,attrib)	dv_clr_region	VB_CLR_REGION
void cmove(y,x)	dv_ctoyx	VB_CTOYX
void curpos(int *yp,int *xp);	dv_getyx	vb_getyx
void doscroll(l,r,t,b,a,at)	dv_scroll	VB_SCROLL
void freescr(SBUF *p);	dv_freesbuf	vb_freesbuf
int incha()	dv_incha	VB_INCHA
int inchar()	dv_incha	VB_INCHA
void outc(c, attrib)	dv_putc	vb_putc
void replace(c);	dv_replace	VB_REPLACE
SBUF *restore(SBUF *b);	dv_restore	vb_restore
SBUF *savescr(l,r,t,b)	dv_save	vb_save

† Unspecified argument types are **int**.

The glue functions are declared in *glue.c*, Listing A.68. One of three sets of functions are compiled depending on the existence of various macros as follows:

One of three sets of glue functions can be compiled, R, V, A.

If this macro is #defined	Use this interface
R	Video-BIOS.
V	Direct-video.
A	Select automatically at run time (autoselect).

You can compile *glue.c* three times, once with R defined, once with V defined, and once with neither defined. The version of *glue.obj* that you link then determines which set of I/O functions are called into the final program.

Select between direct video and video BIOS at run time, autoselect mode.

Initializing the low-level I/O functions, *init()*.

The VIDEO environment.

In “autoselect” mode, the glue functions select one or the other of the sets of I/O functions at run time. The *init()* function is used for this purpose. Normally, when R or V are **#defined** at compile time, *init()* just initializes for the required mode, returning zero if the BIOS is used (if R is defined) or 1 if direct-video functions are used.

If A is **#defined** at compile time, the *init()* function selects either the BIOS or direct-video routines at run time. An *init(0)* call forces use of the BIOS, *init(1)* forces use of direct-video functions, *init(-1)* causes *init()* to examine the *VIDEO* environment variable and choose according to its contents—if *VIDEO* is set to "DIRECT" or "direct", the direct-video functions are used, otherwise the video-BIOS is used. *init()* returns 0 if the BIOS is selected, 1 if the direct-video functions are selected. The one disadvantage to autoselect mode is that both libraries must be in memory at run time, even though you’re using only one of them.

Listing A.68. *glue.c*—Glue Functions for Low-Level I/O

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <tools/debug.h>
4 #include <tools/termlib.h>
5 #include <tools/vbios.h>
6
7 /* GLUE.C      This file glues the curses package to either the video-
8 *               BIOS functions or to the direct-video functions.
9 */
10
11 #if (!defined(R) && !defined(V))
12 #define AUTOSELECT
13#endif
14
15 /***** R ****/
16 #ifdef R
17
18 init(how) { return 0; }
19 cmove(y,x) { return VB_CTOYX (y,x); }
20 curpos(yp,xp) int *yp,*xp; { return vb_getyx (yp,xp); }
21 replace(c) { return VB_REPLACE (c); }
22 doscroll(l,r,t,b,a,at) { return VB_SCROLL (l,r,t,b,a, at ); }
23 inchar() { return VB_INCHA ( ) & 0xff; }
24 incha() { return VB_INCHA ( ); }
25 outc(c, attrib) { return vb_putc (c, attrib); }
26 SBUF *savescr(l,r,t,b) { return vb_save (l,r,t,b); }
27 SBUF *restore(b) SBUF *b; { return vb_restore (b); }
28 freescr(p) SBUF *p; { return vb_freesbuf(p); }
29 clr_region(l,r,t,b,attrib) { return VB_CLR_REGION(l,r,t,b,attrib); }
30
31 int is_direct(){ return 0; }
32#endif
33

```

Listing A.68. continued...

```

34  /*-----*/
35  #ifdef V
36
37  cmove(y,x)          { return dv_ctoyx  (y,x); }
38  curpos(yp,xp) int *yp,*xp; { return dv_getyx  (yp,xp); }
39  replace(c)           { return dv_replace (c); }
40  doscroll(l,r,t,b,a,at) { return dv_scroll (l,r,t,b,a,at); }
41  inchar()             { return dv_incha  (' ') & 0xff; }
42  incha()              { return dv_incha  (''); }
43  outc(c,attrib)       { return dv_putc   (c, attrib); }
44  SBUF *savescr(l,r,t,b) { return dv_save   (l,r,t,b); }
45  SBUF *restore(b) SBUF *b; { return dv_restore (b); }
46  freescr(p)           SBUF *p; { return dv_freesbuf(p); }
47  clr_region(l,r,t,b,attrib) { return dv_clr_region(l,r,t,b,attrib); }
48
49  init( how )           /* Initialize */
50  {
51      if( !dv_init() )
52      {
53          fprintf(stderr, "MGA or CGA in 80-column text mode required\n");
54          exit( 1 );
55      }
56      return 1;
57  }
58
59  int is_direct(){ return 1; };
60  #endif
61
62  /*-----*/
63  #ifdef A
64
65  static int Dv = 0 ;
66
67  init( how )
68  int how;           /* 0=BIOS, 1=direct video, -1=autoselect */
69  {
70      char *p;
71
72      if( Dv == how )
73      {
74          if( how < 0 )
75          {
76              p = getenv( "VIDEO" );
77              Dv = p && ((strcmp(p,"DIRECT")==0 || strcmp(p,"direct")==0));
78          }
79
80          if( Dv && !dv_init() )
81          {
82              fprintf(stderr, "MGA or CGA in 80-column text mode required\n");
83              exit( 1 );
84          }
85      }
86
87      return Dv;
88  }
89
90  /* The following statements all depend on the fact that a subroutine name
91   * (without the trailing argument list and parentheses) evaluates to a temporary
92   * variable of type pointer-to-function. Therefore, function names can be treated
93   * as a pointer to a function in the conditionals, below. For example, curpos() */

```

Listing A.68. continued...

```

94  * calls dv_getyx if Dv is true, otherwise vb_getyx is called. ANSI actually
95  * says that the star isn't necessary. The following should be legal:
96  *
97  *      return (Dv ? dv_ctyx : VB_CTOYX)( y, x );
98  *
99  * but many compilers (including Microsoft) don't accept it. Macros, clearly,
100 * must be treated in a more standard fashion.
101 */
102
103 curpos(yp,xp) int *yp,*xp; {return(Dv? *dv_getyx :*vb_getyx ) (yp,xp); }
104 SBUF *savescr(l,r,t,b) {return(Dv? *dv_save :*vb_save ) (l,r,t,b); }
105 SBUF *restore(b) SBUF *b; {return(Dv? *dv_restore :*vb_restore ) (b); }
106 freescr(p) SBUF *p; {return(Dv? *dv_freesbuf:*vb_freesbuf) (p); }
107 outc(c, attrib) {return(Dv? *dv_putc :*vb_putc ) (c, attrib); }
108
109 replace(c) { if(Dv) dv_replace(c); else VB_REPLACE(c); }
110 cmove(y,x) { if(Dv) dv_ctoyx(y,x); else VB_CTOYX(y,x); }
111 inchar() { return( Dv? dv_incha() : VB_INCHA() ) & 0xff; }
112 incha() { return( Dv? dv_incha() : VB_INCHA() ); }
113
114 doscroll(l,r,t,b,a,at)
115 {
116     if( Dv ) dv_scroll(l,r,t,b,a,at);
117     else     VB_SCROLL(l,r,t,b,a,at);
118 }
119
120 clr_region(l,r,t,b,attrib)
121 {
122     if( Dv ) dv_clr_region(l,r,t,b,attrib);
123     else     VB_CLR_REGION(l,r,t,b,attrib);
124 }
125
126 int is_direct(){ return Dv; };
127 #endif

```

A.11 Window Management: Curses

The low-level I/O functions are really too low level to be useful—higher-level functions are needed. Since I write a lot of code that has to work in both the MS-DOS and UNIX environments, it seemed sensible to mimic a set of UNIX functions under MS-DOS. Among other things, this approach solves various incompatibility problems between the two systems. For example, the Microsoft, MS-DOS C compiler doesn't have UNIX-compatible `fcntl()` or `ioctl()` functions, it doesn't support the `/dev/tty` device for the console (you have to use `/dev/con` or `CON`), and it doesn't provide any sort of `termcap` or curses-compatible function library. Consequently, it's difficult to port code that does low-level I/O from the Microsoft compiler to UNIX.

Curses.

The most serious omission in the earlier list is the lack of a curses library. For the uninitiated, *curses* is a collection of terminal-independent, low-level I/O functions. The package was written by Ken Arnold, then at U.C. Berkeley. These routines let you do things like move the cursor around on the screen, create and delete windows, write text and seek to specific window-relative cursor positions and so forth. The windows can be overlapping and they support individual wrap-around, scrolling, and so forth. The curses functions can talk to virtually any terminal. They accomplish this feat by using the

Termcap database.

termcap terminal database and interface library developed by Bill Joy. The database is an ASCII file that contains definitions for the various escape sequences needed to get around on specific terminals, and the interface library lets you access these machine-specific escape sequences in a portable and transparent way.

The real curses routines talk to the terminals efficiently over a serial line. They always send the minimum number of characters necessary to modify the current screen. Curses keeps two internal images of the screen. One of these reflects what's actually on the screen, and the other is a scratch space that you modify using the various curses functions. When you tell curses to do a refresh, it compares the scratch buffer with the actual screen image and then sends out the minimum number of characters necessary to get these images to match. This behavior is especially important when you're running a program via a modem and characters are coming at 1200 baud. Redrawing the entire screen every time you scroll a four-line by 10-character window is just unacceptable behavior. It takes too long. Curses solves the problem by redrawing only those parts of the screen that have actually changed.

The UNIX curses functions are described in depth in [Arnold] and also in [Haviland]. Since the implementation details can vary somewhat from installation to installation, I strongly recommend that you read your own system documentation as well.

The current section describes the behavior of my own curses implementation, which *occs* uses for window management in interactive debugging environment. My system models Berkeley curses—there are a few minor differences between it and System V versions. I've written several quite complex programs using these functions, programs that maintain several windows on the screen simultaneously, all of which are being updated at different rates. Moreover, the finished programs have ported to UNIX (BSD 4.3) with literally no modification. I have not implemented the entire curses library, however, and I've added a few features to my own package that are not supported by the real curses. You can write a UNIX-compatible program using my functions, but there are minor (but documented) differences that can cause problems if you're not careful. For example, support for overlapping windows is limited (you can't write reliably to a window that has a second window on top of it). Subwindows aren't supported. This last problem means that you can't move a window along with the associated subwindows. Similarly, you have to delete or update the subwindows explicitly, one at a time. Pseudo-subwindows can be used, but only if you don't move them. My package does implement a few handy features not found in the UNIX version, however. You can hide a window without deleting it, and there is a provision for making boxed windows.

This section just describes my own functions. Though I point out the important differences from UNIX, you'll have to read the UNIX documentation too if you intend to write portable code. The two systems do behave differently in many ways. Note that, since these routines are not intended to port to UNIX, I haven't done things like use the `UX()` and `VA_LIST` macros in `<debug.h>` to make them portable.

Curses is designed for efficient serial communication.

Curses portability issues.

A.11.1 Configuration and Compiling

The curses functions talk to the screen using the low-level I/O functions described in the previous section. Curses automatically selects between the CGA and MGA adapters by using the autoselect version of the glue functions. (The CGA must be running in an 80-column mode, though.) Curses normally uses the ROM-BIOS routines, which are slow but portable. If, however, a `VIDEO` environment variable exists, and that variable is set to the string `DIRECT`, the direct-video functions will be used. (Do this with a `set VIDEO=DIRECT` from the DOS prompt.) The BIOS functions are noticeably slower than the direct-video functions. The speed problem is most evident when you are saving and restoring the area under a window. Moving a visible window in an

incremental fashion (one space or line at a time) is a particularly painful process when you're using the BIOS. It's better, in this instance, to hide the window, move it where you want it, and then redisplay it.

curses.lib, *termlib.lib*.

The curses functions are all in a library called *curses.lib*, and the I/O routines are in *termlib.lib*. Both of these must be linked to the final program.

A.11.2 Using Curses

<curses.h>

Curses itself is part macro and part subroutines. The file *<curses.h>* should be **#included** at the top of every file that uses the curses functions. Supported functions are described in this section, grouped functionally.

A.11.2.1 Initialization Functions.

Initialize,
terminate curses.

```
void initscr(void);
void endwin (void);
```

initscr() initializes the curses package. It should be called at the head of your *main()* subroutine, before any other curses functions are called. *endwin()* cleans up. It should always be called before your program exits.

In UNIX programs, the terminal can be left in an unknown state if you abort your program with a BREAK. If you exit abnormally from a program that uses curses, only to find your terminal acting funny (not echoing, not handling tabs or newlines properly, and so forth), you can usually correct the problem by typing *tset* with no arguments. If that doesn't work, try <NL>*reset*<NL> where <NL> is a newline or Ctrl-J. If that doesn't work try *stty cooked echo nl*, and if that doesn't work, hang up and log on again. To avoid this sort of flailing around, it's much better for your program to trap SIGINT and call *endwin()* from within the service routine. Use the following:

```
#include <signal.h>
#include <curses.h>

onintr()
{
    endwin();
    exit( STATUS );
}
main()
{
    signal ( SIGINT, onintr );
    initscr ( void );
}
```

A.11.2.2 Configuration Functions. Once the curses package is initialized, you should determine how your terminal is going to respond to typed characters. Six routines are provided for this purpose:

Set unbuffered,
buffered input.

```
int crmode (void);
int nocrmode(void);
```

These routines control input buffering; *crmode()* disables buffering—characters will be available as soon as they're typed. A *nocrmode()* call cancels a previous *crmode()*. In *nocr* mode, an entire line is read before the first character is returned. The only editing character available in *nocrmode()* is a backspace, which deletes the character to the left of the cursor. Many curses programs use *crmode()*, but some implementations won't work properly unless *nocrmode()* is active.

```
int echo (void);
int noecho(void);
```

Echo characters.
Do not echo.

If `echo()` is called, characters are echoed as they're typed; `noecho()` suppresses the echoing—you'll have to echo the character yourself every time you read a character. The real curses gets very confused when `echo()` is enabled. The problem here is that curses doesn't know about any character that it has not itself written to the screen. Since characters are echoed by the operating system rather than curses, the package doesn't know that they're there. As a consequence, when curses does a screen refresh, it won't delete the characters that it doesn't know about and the screen rapidly fills with unwanted and unerasable characters. It's best always to call `noecho()` at the top of your program. Another echo-related problem found in the MS-DOS versions of curses. Character echo cannot be suppressed with the MS-DOS buffered-input function. So `echo()` and `noecho()` have no effect on a program if `nocrmode()` is active.

```
int nl (void);
int nonl(void);
```

Map newlines.
Do not map newlines.

A `nl()` call causes a newlines ('`\n`') to be converted to a carriage-return, line-feed sequence on output; an input carriage return ('`\r`') is also mapped to a newline. If `nonl()` is called, no mapping is done. It's usually convenient to call `nl()` at the top of your program, but again, many UNIX curses packages fail to work properly unless `nonl()` is specified. You have to do all the '`\n`' to `\r\n` mapping yourself in `nonl` mode, of course.

A.11.2.3 Creating and Deleting Windows. This section describes the functions that create and delete windows. There are two kinds of windows. A normal window and a *subwindow*. In UNIX curses, a subwindow is affected by all commands that also affect the parent. When a parent window is refreshed by curses, all subwindows are refreshed too. Similarly, when you delete or move a parent window, all the subwindows are deleted or moved. This feature is not supported in the current implementation, though you can pretend that subwindows exist if you don't move them. The mechanics of this process are discussed below.

Subwindows.

A default window, called `stdscr`, which occupies the entire screen, is created automatically by `initscr()`, and several functions are provided to modify this window. A `stdscr` variable is defined in `<curses.h>`, and it can be passed as a `WINDOW` pointer to any of the curses functions in a manner analogous to `stdout`. You shouldn't declare `stdscr` explicitly in your program, just use it. For convenience, most functions have a version that uses `stdscr`, rather than an explicit window, in a manner analogous to `putchar()`. If you are using curses only for cursor movement and are not creating additional windows, you can use `stdscr` directly and don't need to use any of the functions described in this section.

The `stdscr` window.

In UNIX applications, it's often convenient to declare all windows as subwindows to `stdscr`. For example, the save-screen mechanism used in `yydebug.c` won't work unless all the windows are subwindows of `stdscr` because it reads characters from `stdscr` itself.

The `stdscr` variable in `<curses.h>`.

```
WINDOW *newwin(int lines, int cols, int begin_y, int begin_x);
WINDOW *subwin(WINDOW *win, int lines, int cols, int begin_y,
               int begin_x);
```

Create window, subwindow.

This function creates a new window, `lines` rows high and `cols` columns wide with the upper-left corner at (`begin_y`, `begin_x`). [All coordinates here are (`y,x`), where `y` is the row number and `x` is the column number. The upper-left

Coordinates are (`y,x`),
`y=`row, `x=`column. Upper
left corner is (0,0).

corner of the screen is (0,0).] A pointer to a WINDOW structure, declared in *<curses.h>*, is returned in a manner analogous to `fopen()` returning a FILE.

Windows are created as visible, unboxed, with scrolling disabled, and with line wrap enabled. Characters that go past the end of line show up at the left edge of the next line, but the window won't scroll when you get to the bottom. The text under the window is saved by default and is restored when the window is moved or deleted. The window is automatically cleared as part of the creation process and there's no way to disable this clearing. All of these defaults can be changed with subroutine calls discussed below.

The `subwin()` call is provided for UNIX compatibility. It is mapped to a `newwin()` call (the first argument is ignored).

Enable or disable scrolling.

```
scrolllok(WINDOW *win, int flag);
```

This macro is passed a WINDOW pointer and a flag. If flag is true, the indicated window is allowed to scroll. Otherwise the window does not scroll and characters that go off the bottom of the window are discarded. Scrolling is always enabled on the `stdscr` window.

Nonstandard: enable or disable line wrap.

```
wrapok(WINDOW *win, int flag);
```

This macro enables or disables line wrap if flag is false, enables it otherwise. Line wrap is enabled by default. When wrapping is disabled, characters written past the edge of the window are discarded rather than appearing on the next line.

Nonstandard:

Box windows.
No boxes.

```
void boxed(void);
void unboxed(void);
```

The lack of support for subwindows complicates management of windows with boxes around them considerably. A UNIX-compatible method for creating boxed windows is described below in the discussion of the `box()` function. It's not practical to use this function if windows will be moved, however. I've solved the problem to some extent by adding a mechanism for creating boxed windows in which the box is an integral part of the window and cannot be overwritten by the text. All windows created with `newwin()` after `boxed()` has been called will have an integral box as part of the window. You can go back to normal, unboxed windows by calling `unboxed()`. IBM box-drawing characters are used for the border.

Nonstandard: change colors.

```
void ground (WINDOW *win, int fore, int back);
void def_ground(int fore, int back);
```

`ground()` lets you change foreground (fore) and background (back) colors on the IBM/PC. The color codes are defined symbolically in *<tools/termlib.h>*. `def_ground()` changes the default foreground and background colors. All windows created after a `def_ground()` call have the indicated colors. Use `ground()` to selectively change the attribute associated with specific characters written to a window; only those characters written to the window after the `ground()` call are affected. If you want to change an existing window's color without modifying the window's contents, you must change the color with `ground()` and then read each character individually and write it back using `winch()` and `addch()`.

BUGS: `delwin()` doesn't know about color changes, you must set the ground back to NORMAL and call `wclear()` before deleting the window.

```
void save (void) ;
void nosave(void) ;
```

Normally, when you create a window, the text under that window is saved so that it can be replaced when the window is moved or deleted. This is a needless waste of memory if you're not going to delete or move the window. Once `nosave()` has been called, the underlying text is not saved by subsequent `newwin()` calls. Saving can be reenabled by calling `save()`. Text under `stdscr` is never saved.

```
delwin(WINDOW *win);
```

Delete a window.

This function removes a window from the screen, usually restoring the text that was underneath the window. If `nosave()` was active when the window was created, however, an empty box is left on the screen when the window is deleted. Overlapping windows must be deleted in the opposite order from which they were created. Similarly, if you move one window on top of a second one, you must delete the top window first. I'd suggest keeping a stack of active windows so that you know the order in which they are created.

A.11.2.4 Subroutines That Affect Entire Windows.

```
WINDOW *hidewin(WINDOW *win);
WINDOW *showwin(WINDOW *win);
```

Nonstandard: hide and redisplay window.

The only way to get rid of a UNIX-curses window is to delete it. This is inconvenient when you really just want to make it invisible for a while and then redisplay it later. The nonstandard `hidewin()` function hides the window (makes it disappear but does not delete it). It returns `NULL` if the window can't be hidden for some reason (usually not enough memory is available to save the image), or the `win` pointer on success. The window can be resurrected sometime later with a `showwin()` call. `showwin()` returns `NULL` (and does nothing) if the window wasn't hidden, otherwise the `win` argument is returned. `hidewin()` always returns its argument—it terminates the program with an error message if it can't get memory to hide the window.

It's okay to move a hidden window, it will just reappear at the new location when you call `showwin()`. You cannot, however, write to a hidden window or the written characters will be lost. Similarly, the restrictions of deleting windows in the opposite order from which they were created also applies to hiding and redisplaying them. When windows can overlap, you must always hide the most-recently displayed window first.

```
int mvwin(WINDOW *win, int y, int x);
```

Move window, absolute.

This function moves a window to an absolute location on the screen. The area under the window is restored, the window is moved to the indicated location, and it is then redisplayed, saving the text that is now under the window first. You should move only the topmost of several overlapping windows. The normal UNIX `mvwin()` returns `ERR` and does nothing if the new coordinates would have moved any part of the window off the screen. I've changed this behavior somewhat by always moving the window as far as I can. The window still won't move off the screen, but it may move a little if it wasn't already at the screen's edge. If you prefer UNIX compatibility, change the 0 in `#define LIKE_UNIX 0` in `mvwin.c` (below) to 1 and recompile.

```
mvwinr(WINDOW *win, int y, int x);
```

Nonstandard: move window, relative.

This nonstandard macro lets you move a window relative to the current window

position. Positive values of *y* move down, negative values move up. Similarly, positive values of *x* move right and negative values go left. For example,

```
mvwinr( win, -5, 10 );
```

moves the window pointed to by *win* five spaces up from, and 10 spaces to the right of its current position. Like `mvwin()`, you can't move the window off the screen.

Refresh `stdscr`.
Refresh window.

```
void refresh (void);
void wrefresh(WINDOW *win);
```

These macros are used by the real curses to do a screen refresh. They force the screen to coincide with the internal representation of the screen. No characters are actually written out to the terminal until a refresh occurs. My own curses writes to the screen immediately, so both of these macros expand to null strings—they are ignored. You'll need them to be able to port code to UNIX, however. `refresh()` refreshes the whole screen (the `stdscr` window and all subwindows of `stdscr`); `wrefresh(win)` is passed a `WINDOW` pointer and refreshes only the indicated window.

Draw box in window.

```
int box(WINDOW *win, int vert, int horiz);
```

This subroutine draws a box in the outermost characters of the window using *vert* for the vertical characters and *horiz* for the horizontal ones. I've extended this function to support the IBM box-drawing characters. If IBM box-drawing characters are specified for *vert* and *horiz*, `box()` will use the correct box-drawing characters for the corners. The box-drawing characters are defined in `<tools/box.h>` as follows:

symbolic value	numeric value	description
HORIZ	0xc4	Single horizontal line.
D_HORIZ	0xcd	Double horizontal line.
VERT	0xb3	Single vertical line.
D_VERT	0xba	Double vertical line.

Boxes may have double horizontal lines and single vertical ones, or vice versa. The UNIX `box()` function uses the vertical character for the corners.

Note that `box()` doesn't draw a box around the window; rather, it draws the box in the outermost characters of the window itself. This means that you can overwrite the border if your output lines are too wide. When you scroll the window, the box scrolls too. Normally, this problem is avoided by using the subwindow mechanism. A large outer window is created and boxed, a smaller subwindow is then created for the text region.

Creating a boxed window
with `box()`.

A function that creates a bordered window in this way is shown in Listing A.69. The outer window is created on line 14, the inner, nested one on line 22. The outer window just holds the box and the inner window is used for characters. This way, you can overflow the inner window and not affect the outer one (that holds the border). The function returns a pointer to the text window. If you plan to move or delete the window, you'll need the pointer to the outer window too, and will have to modify the subroutine accordingly.

Listing A.69. Creating a Boxed Window

```

1  WINDOW *boxwin( lines, cols, y_start, x_start, title )
2  char      *title;
3  {
4      /* This routine works much like the newwin() except that the window has a
5      * box around it that won't be destroyed by writes to the window. It
6      * accomplishes this feat by creating two windows, one inside the other,
7      * with a box drawn in the outer one. It prints the title centered on the
8      * top line of the box. Note that I'm making all windows subwindows of the
9      * stdscr to facilitate the print-screen command.
10     */
11
12     WINDOW *outer, *inner;
13
14     outer = subwin(stdscr, lines, cols, y_start, x_start);
15
16     box( outer, '|', '-' );
17
18     wmove( outer, 0, (cols - strlen(title))/2 );
19     wprintw( outer, "%s", title );
20
21     wrefresh( outer );
22     return subwin( stdscr, lines-2, cols-2, y_start+1, x_start+1 );
23 }
```

A.11.2.5 Cursor Movement and Character I/O.

int move (y, x);	Move cursor.
int wmove(WINDOW *win, y, x);	
void getyx(WINDOW *win, y, x);	Get cursor position (MACRO).

These functions support cursor movement. `move()` moves the cursor to the indicated absolute position on the screen. The upper-left corner of the screen is (0,0). `wmove()` moves the cursor to the relative position within a specific window (pointed to by `win`). The upper-left corner of the window is (0,0). If you try to move past the edge of the window, the cursor will be positioned at the closest edge. The `getyx()` macro loads the current cursor position for a specific window into `y` and `x`. Note that this is a macro, not a subroutine, so you should not precede `y` or `x` with an address-of operator (`&`). A `getyx(stdscr,y,x)` call loads the current absolute cursor position into `y` and `x`.

int getch (void);	Get character and echo to window
int wgetch(WINDOW *win);	

These functions are used for direct keyboard input. `getch()` gets a character from the keyboard and echoes it to `stdscr`, `wgetch()` echoes the character to the indicated window (if `echo()` is enabled, that is). Note that `crmode()` has to be enabled to get the character as soon as it's typed. Otherwise the entire line will be buffered. It's unfortunate that many compiler manufacturers, Microsoft included, have chosen to use `getch()` as the name of their standard direct keyboard-input function. You'll need to specify a `/NOE` switch to Microsoft link to prevent the linker from attempting to call in both versions of the subroutine.

Use Microsoft Link's
`/NOE` switch to avoid
name conflicts.

Read functions.

```
inch  (void);
winch (WINDOW *win);
mvinch (int y, int x);
mvwinch(WINDOW *win, int y, int x);
```

These functions let you read back characters that are displayed on the screen. `inch()` returns the character at the current cursor position (that is, from `stdscr`). `mvinch(y,x)` moves the cursor to the indicated position and then returns the character at that position; the `winch()` and `mvwinch(win,y,x)` versions do the same, but the cursor position is relative to the origin of the specified window. Note that the `y` and `x` coordinates are relative to the origin of the indicated window, not to the entire screen. (0,0) is the top left corner of the window. Some older versions of curses don't support the `mv` versions of this command. Note that all four functions work as if all windows were declared as subwindows to `stdscr`. This is not the case in UNIX systems, in which a read from `stdscr` can get a character from the standard screen that is currently obscured by an overlapping window.

`printf()` to `stdscr`.
`printf()` to window.

```
voidprintw(char *fmt,...);
intwprintw(WINDOW *win, char *fmt,...);
```

These functions are used for formatted output. `printw()` works just like `printf()`; `wprintw()` is the same but it prints to the indicated window, moving the cursor to the correct position in the new window if necessary. (The cursor is moved to the position immediately following the character most recently written to the indicated window.) `printw()` pays no attention to window boundaries, but `wprintw()` wraps when you get to the right edge of the window and the window scrolls when you go past the bottom line (provided that `scrollok()` has been called for the current window).

Character-output functions.

```
voidaddch (int c);
intwaddch (WINDOW *win, int c);
voidaddstr (char *str);
intwaddstr(WINDOW *win, char *str);
```

These functions are the curses equivalents of `putc()`, `putchar()`, `puts()`, and `fputs()`: `addch()` works like `putchar()`; `waddch()` writes a character to the indicated window (and advances the cursor); `addstr()` and `waddstr()` work like `fputs()`, writing a string to `stdscr` or the indicated window. Neither `addstr()` nor `waddstr()` add a newline at the end of the string. `addch()` and `waddch()` treat several characters specially:

- '\n' Clear the line from the current cursor position to the right edge of the window. If `nl()` is active go to the left edge of the next line, otherwise go to the current column on the next line. In addition, if scrolling is enabled, the window scrolls if you're on the bottom line.
- '\t' is expanded into an 8-space field. If the tab goes past the right edge of the window, the cursor wraps to the next line.
- '\r' gets you to the left edge of the window, on the current line.
- '\b' backs up one space but may not back up past the left edge of the window. A nondestructive backspace is used (the character on which the initial cursor sits is not erased). The curses documentation doesn't say that '\b' is handled specially but it does indeed work.
- ESC The ESC character is not handled specially by UNIX but my `waddch()` does do so. Don't use explicit escape sequences if portability is a

consideration. All characters between an ASCII ESC and an alphabetic character (inclusive) are sent to the output but are otherwise ignored. This behavior lets you send escape sequences directly to the terminal to change character attributes and so forth. I'm assuming here that you won't change windows in the middle of an escape sequence.

`waddch()` returns `ERR` (defined in `<curses.h>`) if the character had caused the window to scroll illegally or if you attempt to write to a hidden window.

```
void werase (WINDOW *win);
void erase  (void);
void wclear (WINDOW *win);
void clear   (void);
void wclrtoeol(WINDOW *win);
void clrtoeol (void);
```

Erase window, entire screen, to end of line.

These functions all erase one or more characters. `clear()` and `erase()` both clear the entire screen, `wclear()` and `werase()` both clear only the indicated window, `wclrtoeol()` clears the line from the current cursor position in the indicated window to the right edge of the indicated window, and `clrtoeol()` clears from the current cursor position to the right edge of the screen.

```
int scroll (WINDOW *win);
int wscroll(WINDOW *win, int amt);
```

Scroll window.

These two functions scroll the window: `scroll()` scrolls the indicated window up one line, `wscroll()` scrolls by the indicated amount—up if `amt` is positive, down if it's negative. `wscroll()` is not supported by the UNIX curses. Both functions return 1 if the window scrolled, 0 if not.

There's one caveat about scrolling. The UNIX functions have a bug in them: when a window scrolls, the bottom line is not cleared, leaving a mess on the screen. This problem is not restricted to the `scroll()` subroutine but occurs any time that the window scrolls (as when you write a newline on the bottom line of the window or when a character wraps, causing a scroll). As a consequence, if you're porting to UNIX, you should always do a `clrtoeol()` immediately after either scrolling or printing a newline. Unfortunately, there's no easy way to tell if a window has scrolled because of a character wrap. My curses package doesn't have this problem—the bottom line of the window is always cleared on a scroll.

A.11.3 Curses—Implementation

This section discusses my curses implementation. Note that I haven't bothered to make this code UNIX compatible—I'm assuming that you're going to use the native UNIX curses in that environment. For the most part, the code is straightforward and needs little comment. The `WINDOW` structure is declared on lines one to 17 of `<tools/curses.h>` (Listing A.70). The macros for `bool`, `reg`, `TRUE`, `FALSE`, `ERR`, and `OK` are defined in the UNIX `<curses.h>` file so I've put them here too, though they're not particularly useful. Be careful of:

```
if( foo() == TRUE )
```

The `WINDOW` structure.

Problems with
`if(foo() == TRUE).`

`TRUE` is `#defined` to 1 but, in fact, any nonzero value is true. As a consequence, `foo()` could return a perfectly legitimate true value that didn't happen to be 1, and the test would fail. The test:

```
if( foo() != FALSE )
```

ERR returned by curses
output functions. is safer. Most of the output functions return ERR if a write would normally have caused a scroll but scrolling is disabled.

Listing A.70. *curses.h*—WINDOW Definitions and Macros

```

1 typedef struct _window_
2 {
3     int x_org;           /* X coordinate of upper-left corner */
4     int y_org;           /* Y coordinate of upper-left corner */
5     int x_size;          /* Horizontal size of text area. */
6     int y_size;          /* Vertical size of text area. */
7     int row;             /* Current cursor row (0 to y_size-1) */
8     int col;             /* Current cursor column (0 to x_size-1) */
9     void *image;          /* Image buffer. Holds what used to be
10                            /* under the window. */
11     unsigned wrap_ok :1; /* Line wrap is enabled in this win. */
12     unsigned scroll_ok :1; /* Scrolling permitted in this window */
13     unsigned hidden :1;   /* Window is hidden (nonstandard) */
14     unsigned boxed :1;    /* Window is boxed (nonstandard) */
15     unsigned attrib :8;   /* attribute used for character writes */
16
17 } WINDOW;
18
19 #define bool      unsigned int
20 #define reg        register
21 #define TRUE       (1)
22 #define FALSE      (0)
23 #define ERR        (0)
24 #define OK         (1)
25
26 /*****
27 * The following macros implement many of the curses functions.
28 */
29
30 #define getyx( win, y, x ) \
31             (x) = ((WINDOW*)(win))->col, (y) = ((WINDOW*)(win))->row
32
33 #define refresh()
34 #define scrollok(win,flag) ((win)->scroll_ok = (flag))
35 #define wrapok(win,flag) ((win)->wrap_ok = (flag))
36 #define wrefresh(win) /* empty */
37
38 /*****
39 * Nonstandard Macros: movewin() moves the window relative to the current
40 *                      position. Negative is left or up, positive is right or
41 *                      down. ground() changes the fore and background colors for subsequent writes
42 *                      to the window.
43 */
44
45 #define mvwinr(w,dy,dx) mvwin((w), (w)->y_org - (w)->boxed) + (dy), \
46                                ((w)->x_org - (w)->boxed) + (dx))
47
48 #define ground(win,f,b)   ( win->attrib = ((f) & 0x7f) | ((b) & 0x7f) << 4)
49
50 /*****
51 * Externs for the window functions and #defines to map the standard screen
52 * functions to the stdscr versions. There are a few idiosyncrasies here.
53 * In particular, mvcur() just ignores it's first two arguments and maps to a
54 * move() call. Similarly, subwin() just maps to a newwin() call, and clearok()
```



Listing A.70. continued...

```

55  * isn't supported. You must clear the window explicitly before writing to it.
56  */
57
58 extern WINDOW *stdscr;
59 extern void endwin (void);
60 extern void initscr (void);
61 extern int waddch (WINDOW *, int);
62
63 #define addch(c) waddch(stdscr, c)
64
65 extern int waddstr (WINDOW *, char *);
66 #define addstr(s) waddstr(stdscr,s)
67
68 extern int wclrtoeol (WINDOW * );
69 #define clrtoeol() wclrtoeol(stdscr)
70
71 extern int werase (WINDOW * );
72 #define erase werase(stdscr)
73
74 #define wclear(w) werase(w)
75 #define clear() werase(stdscr)
76
77 extern int wgetch (WINDOW * );
78 #define getch() wgetch( stdscr )
79
80 extern int wmove (WINDOW *, int, int );
81 #define move(y,x) wmove( stdscr,(y),(x) )
82 #define mvcur(oy,ox,y,x) move((y),(x))
83
84 extern int wprintw (WINDOW *, char *, ... );
85 extern int printw (char *, ... );
86
87 extern int wscroll (WINDOW *, int );
88 #define scroll(win) wscroll(win,1)
89
90 extern int winch (WINDOW * );
91 #define inch() winch(stdscr)
92 #define mvinch(y,x) ( wmove(stdscr,y,x), winch(stdscr) )
93 #define mvwinch(w,y,x) ( wmove(w, y,x), winch(w) )
94
95 extern WINDOW *newwin (int ,int ,int ,int );
96 #define subwin(w,a,b,c,d) newwin(a,b,c,d)
97
98 /*****
99  * Externs for functions that don't have stdscr versions
100 */
101
102 extern int box (WINDOW *,int ,int ); /* UNIX functions */
103 extern int crmode (void);
104 extern int delwin (WINDOW * );
105 extern int echo (void);
106 extern int mvwin (WINDOW *win, int y, int x);
107 extern int nl (void);
108 extern int nocrmode (void);
109 extern int noecho (void);
110 extern int nonl (void);
111 extern void boxed ( void ); /* Nonstandard functions */
112 extern void unboxed ( void );
113 extern void save ( void );
114 extern void nosave ( void );

```



Listing A.70. continued...

```
115 extern void def_ground ( int, int );
```

The comma operator in `mvinch()` and `mvwinch()`. Writing multiple-statement macros.

The `mvinch()` and `mvwinch()` macros on lines 92 and 93 of Listing A.70 use the comma, or *sequence*, operator. The comma operator evaluates from left to right, and the entire expression evaluates to the rightmost object in the list. For example, the `mvinch()` looks like:

```
#define mvinch(y,x) (move(y,x), inch())
```

An equivalent subroutine is:

```
mvinch(y,x)
{
    move(y,x);
    return inch();
}
```

The comma operator is used because there are two statements that have to be executed, the `move()` call and the `inch()` call. Were you to define the macro as:

```
#define mvinch(y,x) move(y,x); inch()
```

the following code wouldn't work:

```
if( condition )
    mvinch(y,x);
```

because it would expand to:

```
if( condition )
    move(y,x);
    inch();
```

Putting curly braces around the statements doesn't help. For example:

```
#define mvinch(y,x) {move(y,x); inch();}

if( condition )
    mvinch(y, x);
else
    something();
```

expands to

```
if( condition )
{
    move(y,x);
    inch();
}
;
else
    something();
```

Here the `else` will try to bind with the semicolon, which is a perfectly legitimate statement in C, causing a *No matching if for else* error message. Though the comma operator solves both of these problems, it isn't very readable. I don't recommend using it unless you must. Never use it if curly braces will work in a particular application. The remainder of curses is in Listings A.71 to A.87.

Listing A.71. *box.h*—IBM/PC Box-Drawing Characters

```

1  /*-
2   * BOX.H: #defines for the box-drawing characters
3  */
4  *      The names are:
5  *
6  *      UL      Upper left corner
7  *      UR      Upper right corner
8  *      LL      lower left corner
9  *      LR      lower right corner
10 *     CEN     Center (intersection of two lines)
11 *     TOP     Tee with the flat piece on top
12 *     BOT     Bottom tee
13 *     LEFT    Left tee
14 *     RIGHT   Right tee
15 *     HORIZ   Horizontal line
16 *     VERT    Vertical line.
17 *
18 *           UL      -TOP-      UR          HORIZ
19 *                   |          |
20 *                   L          R          V
21 *                   E          I          E
22 *                   F--      -CEN-      --G          R
23 *                   T          |          H          T
24 *                   |          T
25 *                   |
26 *           LL      -BOT-      LR
27 *
28 * The D_XXX defines have double horizontal and vertical lines.
29 * The HD_XXX defines have double horizontal lines and single vertical lines.
30 * The VD_XXX defines have double vertical lines and single horizontal lines.
31 *
32 * If your terminal is not IBM compatible, #define all of these as '+' (except
33 * for the VERT #defines, which should be a '|', and the HORIZ #defines, which
34 * should be a '-') by #defining NOT_IBM_PC before including this file.
35 */
36
37 #ifdef NOT_IBM_PC
38 #   define IBM_BOX(x)
39 #   define OTHER_BOX(x) x
40 #else
41 #   define IBM_BOX(x)  x
42 #   define OTHER_BOX(x)
43 #endif
44
45 #define VERT          IBM_BOX( 179 )  OTHER_BOX( ' | ' )
46 #define RIGHT         IBM_BOX( 180 )  OTHER_BOX( '+ ' )
47 #define UR            IBM_BOX( 191 )  OTHER_BOX( '+ ' )
48 #define LL            IBM_BOX( 192 )  OTHER_BOX( '+ ' )
49 #define BOT           IBM_BOX( 193 )  OTHER_BOX( '+ ' )
50 #define TOP           IBM_BOX( 194 )  OTHER_BOX( '+ ' )
51 #define LEFT          IBM_BOX( 195 )  OTHER_BOX( '+ ' )
52 #define HORIZ         IBM_BOX( 196 )  OTHER_BOX( '- ' )
53 #define CEN           IBM_BOX( 197 )  OTHER_BOX( '+ ' )
54 #define LR            IBM_BOX( 217 )  OTHER_BOX( '+ ' )
55 #define UL            IBM_BOX( 218 )  OTHER_BOX( '+ ' )
56 #define D_VERT        IBM_BOX( 186 )  OTHER_BOX( ' | ' )
57 #define D_RIGHT       IBM_BOX( 185 )  OTHER_BOX( '+ ' )
58 #define D_UR          IBM_BOX( 187 )  OTHER_BOX( '+ ' )
59 #define D_LL          IBM_BOX( 200 )  OTHER_BOX( '+ ' )

```



Listing A.71. continued...

```

60 #define D_BOT      IBM_BOX( 202 ) OTHER_BOX( '+' )
61 #define D_TOP      IBM_BOX( 203 ) OTHER_BOX( '+' )
62 #define D_LEFT     IBM_BOX( 204 ) OTHER_BOX( '+' )
63 #define D_HORIZ    IBM_BOX( 205 ) OTHER_BOX( '-' )
64 #define D_CEN      IBM_BOX( 206 ) OTHER_BOX( '+' )
65 #define D_LR       IBM_BOX( 188 ) OTHER_BOX( '+' )
66 #define D_UL       IBM_BOX( 201 ) OTHER_BOX( '+' )
67 #define HD_VERT    IBM_BOX( 179 ) OTHER_BOX( '|' )
68 #define HD_RIGHT   IBM_BOX( 181 ) OTHER_BOX( '+' )
69 #define HD_UR      IBM_BOX( 184 ) OTHER_BOX( '+' )
70 #define HD_LL      IBM_BOX( 212 ) OTHER_BOX( '+' )
71 #define HD_BOT     IBM_BOX( 207 ) OTHER_BOX( '+' )
72 #define HD_TOP     IBM_BOX( 209 ) OTHER_BOX( '+' )
73 #define HD_LEFT    IBM_BOX( 198 ) OTHER_BOX( '+' )
74 #define HD_HORIZ   IBM_BOX( 205 ) OTHER_BOX( '-' )
75 #define HD_CEN     IBM_BOX( 216 ) OTHER_BOX( '+' )
76 #define HD_LR      IBM_BOX( 190 ) OTHER_BOX( '+' )
77 #define HD_UL      IBM_BOX( 213 ) OTHER_BOX( '+' )
78 #define VD_VERT    IBM_BOX( 186 ) OTHER_BOX( '|' )
79 #define VD_RIGHT   IBM_BOX( 182 ) OTHER_BOX( '+' )
80 #define VD_UR      IBM_BOX( 183 ) OTHER_BOX( '+' )
81 #define VD_LL      IBM_BOX( 211 ) OTHER_BOX( '+' )
82 #define VD_BOT     IBM_BOX( 208 ) OTHER_BOX( '+' )
83 #define VD_TOP     IBM_BOX( 210 ) OTHER_BOX( '+' )
84 #define VD_LEFT    IBM_BOX( 199 ) OTHER_BOX( '+' )
85 #define VD_HORIZ   IBM_BOX( 196 ) OTHER_BOX( '-' )
86 #define VD_CEN     IBM_BOX( 215 ) OTHER_BOX( '+' )
87 #define VD_LR      IBM_BOX( 189 ) OTHER_BOX( '+' )
88 #define VD_UL      IBM_BOX( 214 ) OTHER_BOX( '+' )

```

Listing A.72. cur.h—Curses Functions #includes

```

1 #include <stdio.h>
2 #include <stdlib.h>
3 #include <ctype.h>
4 #include <curses.h>          /* routines in the library.           */
5 #include <stdarg.h>          /* va_list and va_start (ANSI)      */
6 #include <tools/debug.h>
7 #include <tools/termlib.h>
8 #include <tools/box.h>         /* of IBM box-drawing characters    */
9 #include "proto.h"            /* function prototypes for curses functions */
10                         /* (distributed on disk but not printed here) */

```

Listing A.73. box.c—Draw Box in Window

```

1 #include "cur.h"
2
3 box( win, vert, horiz )
4 WINDOW *win;
5 {
6     /* Draws a box in the outermost characters of the window using vert for
7      * the vertical characters and horiz for the horizontal ones. I've
8      * extended this function to support the IBM box-drawing characters. That
9      * is, if IBM box-drawing characters are specified for vert and horiz,
10     * box() will use the correct box-drawing characters in the corners. These

```

Listing A.73. continued...

```

11      * are defined in * box.h as:
12      *
13      *      HORIZ    (0xc4)  single horizontal line
14      *      D_HORIZ  (0xcd)  double horizontal line.
15      *      VERT     (0xb3)  single vertical line
16      *      D_VERT   (0xba)  double vertical line.
17      */
18
19      int      i, nrows;
20      int      ul, ur, ll, lr;
21      int      oscroll, nowrap, oy, ox;
22
23      getyx( win, oy, ox );
24      oscroll      = win->scroll_ok; /* Disable scrolling and line wrap      */
25      nowrap       = win->wrap_ok;   /* in case window uses the whole screen */
26      win->scroll_ok = 0;
27      win->wrap_ok  = 0;
28
29      if( !((horiz==HORIZ || horiz==D_HORIZ) && (vert ==VERT || vert ==D_VERT)) )
30          ul = ur = ll = lr = vert ;
31      else
32      {
33          if( vert == VERT )
34          {
35              if(horiz==HORIZ)
36                  ul=UL,      ur=UR,      ll=LL,      lr=LR;
37              else
38                  ul=HD_UL,  ur=HD_UR,  ll=HD_LL,  lr=HD_LR;
39          }
40          else
41          {
42              if( horiz == HORIZ )
43                  ul=VD_UL,  ur=VD_UR,  ll=VD_LL,  lr=VD_LR;
44              else
45                  ul=D_UL,   ur=D_UR,   ll=D_LL,   lr=D_LR;
46          }
47      }
48
49      wmove ( win, 0, 0 );
50      waddch( win, ul );                      /* Draw the top line */
51
52      for( i = win->x_size-2; --i >= 0 ; )
53          waddch( win, horiz );
54
55      waddch( win, ur );
56      nrows = win->y_size - 2 ;                /* Draw the two sides */
57
58      i = 1 ;
59      while( --nrows >= 0 )
60      {
61          wmove ( win, i, 0 );
62          waddch( win, vert );
63          wmove ( win, i++, win->x_size - 1 );
64          waddch( win, vert );
65      }
66
67      wmove ( win, i, 0 );                      /* Draw the bottom line */
68      waddch( win, ll );
69

```



Listing A.73. continued...

```

70     for( i = win->x_size-2; --i >= 0 ; )
71         waddch( win, horiz );
72
73     waddch( win, lr );
74     wmove ( win, oy, ox );
75     win->scroll_ok = oscroll ;
76     win->wrap_ok   = owrap ;
77 }
```

Listing A.74. delwin.c—Delete Window

```

1 #include "cur.h"
2
3 delwin( win )
4 WINDOW *win;
5 {
6     /* Copy the saved image back onto the screen and free the memory used for
7      * the buffer.
8      */
9
10    if( win->image )
11    {
12        restore ( (SBUF *) win->image );
13        freescr ( (SBUF *) win->image );
14    }
15    free( win );
16 }
```

Listing A.75. hidewin.c—Hide a Window

```

1 #include "cur.h"
2
3 WINDOW *hidewin( win )
4 WINDOW *win;
5 {
6     /* Hide a window. Return NULL and do nothing if the image wasn't saved
7      * originally or if it's already hidden, otherwise hide the window and
8      * return the win argument. You may not write to a hidden window.
9      */
10
11    SBUF *image;
12
13    if( !win->image || win->hidden )
14        return NULL;
15
16    image = savescr(((SBUF*)(win->image))->left,((SBUF*)(win->image))->right,
17                  ((SBUF*)(win->image))->top, ((SBUF*)(win->image))->bottom );
18
19    restore( (SBUF *) win->image );
20    freescr( (SBUF *) win->image );
21    win->image = image;
22    win->hidden = 1;
23    return( win );
24 }
```

Listing A.76. *initscr.c*— Initialize Curses

```

1 #include "cur.h"
2
3 WINDOW *stdscr;
4
5 void endwin() /* Clean up as required */
6 {
7     cmove( 24, 0 );
8 }
9
10 void initscr()
11 {
12     /* Creates stdscr. If you want a boxed screen, call boxed() before this
13      * routine. The underlying text is NOT saved. Note that the atexit call
14      * insures a clean up on normal exit, but not with a Ctrl-Break, you'll
15      * have to call signal() to do that.
16     */
17
18     nosave();
19     init(-1);
20     stdscr = newwin( 25, 80, 0, 0 );
21     save();
22     atexit( endwin );
23 }
```

Listing A.77. *mvwin.c*— Move a Window

```

1 #include "cur.h"
2
3 /* Move a window to new absolute position. This routine will behave in one of
4  * two ways, depending on the value of LIKE_UNIX when the file was compiled.
5  * If the #define has a true value, then mvwin() returns ERR and does nothing
6  * if the new coordinates would move the window past the edge of the screen.
7  * If LIKE_UNIX is false, ERR is still returned but the window is moved flush
8  * with the right edge if it's not already there. ERR says that the window is
9  * now flush with the edge of the screen. In both instances, negative
10 * coordinates are silently set to 0.
11 */
12
13 #define LIKE_UNIX 0
14
15 #if ( LIKE_UNIX )
16 #    define UNIX(x) x
17 #    define DOS(x)
18 #else
19 #    define UNIX(x)
20 #    define DOS(x) x
21#endif
22
23 mvwin( win, y, x )
24 WINDOW *win;
25 {
26     int old_x, old_y, xsize, ysize, delta_x, delta_y, visible;
27     SBUF *image;
28
29     if( win == stdscr ) /* Can't move stdscr without it going */
30         /* off the screen. */ */
31 }
```

Listing A.77. continued...

```

32     /* Get the actual dimensions of the window: compensate for a border if the
33      * window is boxed.
34      */
35
36     old_x = win->x_org - win->boxed ;
37     old_y = win->y_org - win->boxed ;
38     xsize = win->x_size + (win->boxed * 2);
39     ysize = win->y_size + (win->boxed * 2);
40
41     /* Constrain x and y so that the window can't go off the screen */
42
43     x = max( 0, x );
44     y = max( 0, y );
45     if( x + xsize > 80 )
46     {
47         UNIX( return ERR;           )
48         DOS ( x = 80 - xsize;    )
49     }
50     if( y + ysize > 25 )
51     {
52         UNIX( return ERR;           )
53         DOS ( y = 25 - ysize;    )
54     }
55
56     delta_x = x - old_x;          /* Adjust coordinates. */
57     delta_y = y - old_y;
58
59     if( delta_y == 0 && delta_x == 0 )
60         return ERR;
61
62     if( visible = !win->hidden )
63         hidewin( win );
64
65     win->y_org    += delta_y;
66     win->x_org    += delta_x;
67     image          = (SBUF *) win->image;
68     image->top     += delta_y;
69     image->bottom  += delta_y;
70     image->left    += delta_x;
71     image->right   += delta_x;
72
73     if( visible )
74         showwin( win );
75     return( OK );
76 }
```

Listing A.78. showwin.c—Display a Previously Hidden Window

```

1 #include "cur.h"
2
3 WINDOW *showwin( win )
4 WINDOW *win;
5 {
6     /* Make a previously hidden window visible again. Return NULL and do
7      * nothing if the window wasn't hidden, otherwise return the win argument.
8      */
9 }
```

Listing A.78. continued...

```

10    SBUF *image;
11
12    if( !win->hidden )
13        return( NULL );
14
15    image = savescr(((SBUF*)(win->image))->left, ((SBUF*)(win->image))->right,
16                                ((SBUF*)(win->image))->top,   ((SBUF*)(win->image))->bottom);
17
18    restore( (SBUF *) win->image );
19    freescr( (SBUF *) win->image );
20    win->image = image;
21    win->hidden = 0;
22
23    /* Move the cursor to compensate for windows that were moved while they
24     * were hidden.
25     */
26    cmove( win->y_org + win->row, win->x_org + win->col );
27    return( win );
28 }
```

Listing A.79. waddstr.c—Write String to Window

```

1 #include "cur.h"
2
3 waddstr( win, str )
4 WINDOW *win;
5 char *str;
6 {
7     while( *str )
8         waddch( win, *str++ );
9 }
```

Listing A.80. wclrtoeo.c—Clear from Cursor to Edge of Window

```

1 #include "cur.h"
2
3 wclrtoeo( win )
4 WINDOW *win;
5 {
6     /* Clear from cursor to end of line, the cursor isn't moved. The main
7      * reason that this is included here is because you have to call it after
8      * printing every newline in order to compensate for a bug in the real
9      * curses. This bug has been corrected in my curses, so you don't have to
10     * use this routine if you're not interested in portability.
11     */
12
13     clr_region( win->x_org + win->col , win->x_org + (win->x_size - 1),
14                 win->y_org + win->row , win->y_org + win->row,
15                 win->attrib
16 }
```

Listing A.81. *werase.c*—Erase Entire Window

```

1 #include "cur.h"
2
3 werase( win )
4 WINDOW *win;
5 {
6     clr_region( win->x_org, win->x_org + (win->x_size - 1),
7                 win->y_org, win->y_org + (win->y_size - 1),
8                 win->attrib
9                 );
10    cmove( win->y_org, win->x_org );
11    win->row = 0;
12    win->col = 0;
13 }
```

Listing A.82. *winch.c*—Move Window-Relative Cursor and Read Character

```

1 #include "cur.h"
2
3 int      winch( win )
4 WINDOW *win;
5 {
6     int y, x, c;
7
8     curpos( &y, &x );
9     cmove( win->y_org + win->row, win->x_org + win->col );
10    c = inchar();
11    cmove( y, x );
12    return c;
13 }
```

Listing A.83. *wincreat.c*—Create a Window

```

1 #include "cur.h"
2 #include <tools/box.h>
3
4 /*****
5  * Window creation functions.
6  * Standard Functions:
7  *
8  *      WINDOW      *newwin( lines, cols, begin_y, begin_x )
9  *                  creates a window
10 *
11 * Nonstandard Functions:
12 *
13 *      save()      Area under all new windows is saved (default)
14 *      nosave()    Area under all new windows is not saved
15 *
16 *      boxed()      Window is boxed automatically.
17 *      unboxed()    Window is not boxed (default)
18 *      def_ground(f,b) Set default foreground color to f, and background color
19 *                  to b.
20 *
21 */
22
```

Listing A.83. continued...

```

23 PRIVATE int Save    = 1;          /* Save image when window created      */
24 PRIVATE int Box     = 0;          /* Windows are boxed                  */
25 PRIVATE int Attrib  = NORMAL;    /* Default character attribute byte    */
26
27 void save () { Save = 1; }
28 void nosave () { Save = 0; }
29 void boxed () { Box = 1; }
30 void unboxed () { Box = 0; }
31
32 void def_ground( f, b ) { Attrib = (f & 0x7f) | ((b & 0x7f) << 4); }
33
34 /***** */
35
36 WINDOW *newwin( lines, cols, begin_y, begin_x )
37
38 int     cols;      /* Horizontal size (including border) */
39 int     lines;     /* Vertical size (including border)   */
40 int     begin_y;   /* X coordinate of upper-left corner */
41 int     begin_x;   /* Y coordinate of upper-left corner */
42 {
43     WINDOW     *win;
44
45     if( !(win = (WINDOW *) malloc( sizeof(WINDOW) )) )
46     {
47         fprintf(stderr,"Internal error (newwin): Out of memory\n");
48         exit(1);
49     }
50
51     if( cols > 80 )
52     {
53         cols    = 80;
54         begin_x = 0;
55     }
56     else if( begin_x + cols > 80 )
57         begin_x = 80 - cols;
58
59     if( lines > 25 )
60     {
61         lines   = 25;
62         begin_y = 0;
63     }
64     else if( begin_y + lines > 25 )
65         begin_x = 25 - cols;
66
67     win->x_org      = begin_x ;
68     win->y_org      = begin_y ;
69     win->x_size     = cols   ;
70     win->y_size     = lines   ;
71     win->row        = 0 ;
72     win->col        = 0 ;
73     win->scroll_ok  = 0 ;
74     win->wrap_ok    = 1 ;
75     win->boxed      = 0 ;
76     win->hidden     = 0 ;
77     win->attrib     = Attrib ;
78     win->image      = !Save ? NULL : savescr( begin_x, begin_x + (cols - 1) ,
79                                         begin_y, begin_y + (lines - 1) );
80     werase(win);
81

```



Listing A.83. continued...

```

82     if( Box )                      /* Must be done last */
83     {
84         box( win, VERT, HORIZ );    /* Box it first           */
85         win->boxed   = 1;
86         win->x_size -= 2;          /* Then reduce window size */
87         win->y_size -= 2;          /* so that the box won't */
88         win->x_org  += 1;          /* be overwritten.        */
89         win->y_org  += 1;
90
91         cmove( win->y_org, win->x_org );
92     }
93     return win;
94 }
```

Listing A.84. winio.c—Miscellaneous Window I/O Functions

```

1 #include "cur.h"
2
3 /-----
4 * WINIO.C      Lowest level I/O routines:
5 *      waddch(win,c)      wgetch(win)
6 *      echo()            noecho()
7 *      nl()              nonl()
8 *      crmode()          nocrmode()
9 *-----
10 */
11
12 PRIVATE int Echo    = 1;          /* Echo enabled           */
13 PRIVATE int Crmode = 0;          /* If 1, use buffered input */
14 PRIVATE int Nl     = 1;          /* If 1, map \r to \n on input */
15                         /* and map both to \n\r on output */
16 echo    () { Echo = 1; }
17 noecho   () { Echo = 0; }
18 nl      () { Nl   = 1; }
19 nonl    () { Nl   = 0; }
20
21 crmode()
22 {
23     Crmode = 1;
24 }
25
26 nocrmode()
27 {
28     Crmode = 0;
29 }
30
31 /-----*/
32
33 static char  *getbuf(win, buf)
34 WINDOW  *win;
35 char    *buf;
36 {
37     /* Get a buffer interactively. ^H is a destructive backspace. This routine
38     * is mildly recursive (it's called from wgetch() when Crmode is false.
39     * The newline is put into the buffer. Returns it's second argument.
40     */
41 }
```

Listing A.84. continued...

```

42     register int c;
43     char          *sbuf = buf;
44
45     Crmode = 1;
46     while( (c = wgetch(win)) != '\n' && c != '\r' )
47     {
48         switch( c )
49         {
50             case '\b': if( --buf >= sbuf )
51                         wprintw( win, " \b" );
52                     else
53                     {
54                         wprintw( win, " " );
55                         putchar('\007');
56                         buf = sbuf;
57                     }
58                     break;
59
60             default:  *buf++ = c;
61                     break;
62         }
63     }
64     *buf++ = c      /* Add line terminator (\n or \r) */
65     *buf = '\0';
66     Crmode = 0 ;
67     return sbuf ;
68 }
69
70 /***** */
71
72 int      wgetch( win )
73 WINDOW  *win;
74 {
75     /* Get a character from DOS without echoing. We need to do this in order
76      * to support (echo/noecho). We'll also do noncrmode input buffering here.
77      * Maximum input line length is 132 columns.
78      *
79      * In nocrmode(), DOS functions are used to get a line and all the normal
80      * command-line editing functions are available. Since there's no way to
81      * turn off echo in this case, characters are echoed to the screen
82      * regardless of the status of echo(). In order to retain control of the
83      * window, input fetched for wgetch() is always done in crmode, even if
84      * Crmode isn't set. If nl() mode is enabled, carriage return (Enter, ^M)
85      * and linefeed (^J) are both mapped to '\n', otherwise they are not mapped.
86      *
87      * Characters are returned in an int. The high byte is 0 for normal
88      * characters. Extended characters (like the function keys) are returned
89      * with the high byte set to 0xff and the character code in the low byte.
90      * Extended characters are not echoed.
91      */
92
93     static unsigned char  buf[ 133 ] = { 0 };
94     static unsigned char  *p          = buf;
95     static    int        numchars   = -1;
96     register    int        c;
97
98     if( !Crmode )
99     {

```



Listing A.84. continued...

```

100     if( !*p )
101         p = getbuf( win, buf );
102     return( Nl && *p == '\r' ) ? '\n' : *p++ ;
103 }
104 else
105 {
106     if( (c = bdos(8,0,0) & 0xff) == ('Z'-'@') )      /* Ctrl-Z */
107         return EOF;
108     else if( !c )                                     /* Extended char */
109         c = ~0xff | bdos(8,0,0) ;
110     else
111     {
112         if( c == '\r' && Nl )
113             c = '\n';
114
115         if( Echo )
116             waddch( win, c );
117     }
118     return c;
119 }
120 }
121
122 /*-----*/
123
124 int      waddch( win, c )
125 WINDOW *win;
126 int      c;
127 {
128     /* Print a character to an active (not hidden) window:
129     * The following are handled specially:
130     */
131     /* \n Clear the line from the current cursor position to the right edge
132     * of the window. Then:
133     *     if nl() is active:
134     *         go to the left edge of the next line
135     *     else
136     *         go to the current column on the next line
137     * In addition, if scrolling is enabled, the window scrolls if you're
138     * on the bottom line.
139     * \t is expanded into an 8-space field. If the tab goes past the right
140     * edge of the window, the cursor wraps to the next line.
141     * \r gets you to the beginning of the current line.
142     * \b backs up one space but may not back up past the left edge of the
143     * window. Nondestructive. The curses documentation doesn't say that
144     * \b is handled explicitly but it does indeed work.
145     * ESC This is not standard curses but is useful. It is valid only if
146     * DOESC was true during the compile. All characters between an ASCII
147     * ESC and an alphabetic character are sent to the output but are other
148     * wise ignored. This lets you send escape sequences directly to the
149     * terminal if you like. I'm assuming here that you won't change
150     * windows in the middle of an escape sequence.
151     */
152     /* Return ERR if the character would have caused the window to scroll
153     * illegally, or if you attempt to write to a hidden window.
154     */
155
156     static int      saw_esc = 0;
157     int             rval    = OK;
158

```



Listing A.84. continued...

```
159     if( win->hidden )
160         return ERR;
161
162     cmove(win->y_org + win->row, win->x_org + win->col);
163
164 #ifdef DOESC
165     if( saw_esc )
166     {
167         if( isalpha( c ) )
168             saw_esc = 0;
169         outc(c, win->attrib );
170     }
171     else
172 #endif
173     {
174         switch( c )
175         {
176 #ifdef DOESC
177             case '\033': if( saw_esc )
178                 saw_esc = 0;
179             else
180             {
181                 saw_esc = 1;
182                 outc('\033', win->attrib );
183             }
184             break;
185 #endif
186             case '\b': if( win->col > 0 )
187             {
188                 outc( '\b', win->attrib );
189                 --( win->col );
190             }
191             break;
192
193             case '\t': do {
194                 waddch( win, ' ' );
195             } while( win->col % 8 );
196             break;
197
198             case '\r': win->col = 0;
199             cmove(win->y_org + win->row, win->x_org);
200             break;
201
202             default: if( (win->col + 1) < win->x_size )
203             {
204                 /* If you're not at the right edge of the window, */
205                 /* print the character and advance. */
206
207                 ++win->col;
208                 outc(c, win->attrib );
209                 break;
210             }
211             replace(c);           /* At right edge, don't advance */
212             if( !win->wrap_ok )
213                 break;
214
215             /* otherwise wrap around by falling through to newline */
216 }
```

Listing A.84. continued...

```

217     case '\n': if( c == '\n' )           /* Don't erase character at far */
218                     wclrtoeol( win ); /* right of the screen.          */
219
220         if( NL )
221             win->col = 0;
222
223         if( ++(win->row) >= win->y_size )
224         {
225             rval = wscroll( win, 1 );
226             --( win->row );
227         }
228         cmove( win->y_org + win->row, win->x_org + win->col );
229         break;
230     }
231 }
232 return rval;
233 }
```

Listing A.85. *wmove.c*—Move Window-Relative Cursor

```

1 #include "cur.h"
2
3 wmove( win, y, x )
4 WINDOW *win;
5 {
6     /* Go to window-relative position. You can't go outside the window. */
7
8     cmove( win->y_org + (win->row = min(y,win->y_size-1)) ,
9            win->x_org + (win->col = min(x,win->x_size-1)) );
10 }
```

Listing A.86. *wprintw.c*—Formated Print to Window

```

1 #include "cur.h"
2
3 PRIVATE int Errcode = OK;
4
5 PRIVATE wputc(c, win)
6 int     c;
7 WINDOW *win;
8 {
9     Errcode |= waddch( win, c );
10 }
11
12 wprintw( win, fmt, ... )
13 WINDOW *win;
14 char   *fmt;
15 {
16     va_list args;
17     va_start( args, fmt );
18
19     Errcode = OK;
20     prnt( wputc, win, fmt, args );
21     va_end( args );
22     return Errcode;
```



Listing A.86. continued...

```

23 }
24
25 printw( fmt, ... )
26     *fmt;
27 {
28     va_list args;
29     va_start( args, fmt );
30
31     Errcode = OK;
32     prnt( wputc, stdscr, fmt, args );
33     va_end( args );
34     return Errcode;
35 }
```

Listing A.87. wscroll.c— Scroll Window

```

1 #include "cur.h"
2
3 /*****
4  * Scroll the window if scrolling is enabled. Return 1 if we scrolled. (I'm
5  * not sure if the UNIX function returns 1 on a scroll but it's convenient to
6  * do it here. Don't assume anything about the return value if you're porting
7  * to UNIX. Wscroll() is not a curses function. It lets you specify a scroll
8  * amount and direction (scroll down by -amt if amt is negative); scroll()
9  * is a macro that evaluates to a wscroll call with an amt of 1. Note that the
10 * UNIX curses gets very confused when you scroll explicitly (using scroll()). 
11 * In particular, it doesn't clear the bottom line after a scroll but it thinks
12 * that it has. Therefore, when you try to clear the bottom line, it thinks that
13 * there's nothing there to clear and ignores your wclrtoeol() commands. Same
14 * thing happens when you try to print spaces to the bottom line; it thinks
15 * that spaces are already there and does nothing. You have to fill the bottom
16 * line with non-space characters of some sort, and then erase it.
17 */
18
19 wscroll(win, amt)
20 WINDOW *win;
21 {
22     if( win->scroll_ok )
23         doscroll( win->x_org, win->x_org + (win->x_size-1),
24                    win->y_org, win->y_org + (win->y_size-1), amt, win->attrib );
25
26     return win->scroll_ok ;
27 }
```

Notes on Pascal Compilers

Though the material covered in Chapter Six applies to most compilers, languages such as Pascal, which support nested subroutine declarations, arguments passed by reference, and so forth, have their own set of problems.

B.1 Subroutine Arguments

Passing by value.

The first issue is differentiating between arguments passed by value and arguments passed by reference. With the exception of arrays, all arguments in C are passed by value—the argument’s value (the contents of a variable, for example) is passed to the subroutine rather than the argument itself. Pascal doesn’t differentiate between arrays and other types, however. Unless you tell the compiler to do otherwise, everything is passed by value, so the entire array must be copied onto the stack as part of the subroutine call. Similarly, records must be copied to the stack to pass them to a subroutine (as must structures in ANSI C).

Passing by reference.

In a call by reference, the called subroutine can modify the contents of a variable local to the calling subroutine. A reference to the variable is passed, not its contents. In C, you’d do this by passing a pointer to the variable, and this approach works in Pascal as well, except that the compiler takes care of the details. If a subroutine argument is declared using the **var** keyword, then the address of the object is passed to the subroutine. The subroutine itself must access that object indirectly through the pointer.

An alternative approach is not practical in most languages, but should be mentioned primarily so that you don’t use it. You can pass all arguments to the subroutine by value, just as if they were arguments to a C subroutine. Rather than discarding the arguments after the subroutine returns, however, you can pop their modified values from the stack back into the original variables. This approach has the advantage of simplifying the called subroutine’s life considerably because it doesn’t have to keep track of the extra level of indirection. The obvious disadvantage is that large objects, such as arrays, require a lot of stack space if passed by value. A less obvious, but more serious, disadvantage is illustrated by the following subroutine:

```

function tweedle( var dum: integer; var dee: integer ): integer;
begin
    dum = 1;
    return dum + dee;
end

```

tweedle()'s return value is undefined when a copy-in/copy-out strategy is used and the arguments are identical, as follows:

```

var
    arg:     integer;
    result: integer;
begin
    arg    := 100;
    result := tweedle( arg, arg );
end

```

The assignment dum=1 should modify both dum and dee and the function should return 2. If the arguments are copied onto the stack, however, dum alone is modified, so 101 is returned. The other problem is the undefined order of evaluation in the arguments. There's no way to tell which of the two arguments will be popped first when the subroutine returns. If the left argument, which corresponds to dum, is popped first, then arg is set to 100; otherwise, arg is set to 1.

B.2 Return Values

The next issue is a function's return value. Aggregate objects, such as arrays and records, can't be returned in registers because they're too big. The usual approach is to allocate an area of static memory using `malloc()` or some similar memory allocator, copy the aggregate object into this memory, and return a pointer to the region in a register. You can also reserve a fixed region of static memory for this purpose—all functions copy the aggregate object to the same fixed region of memory. A pointer doesn't have to be returned in this case because everybody knows where the returned object can be found. On the other hand, the size of the fixed region will limit the size of the returned object.

Returning aggregate objects.

B.3 Stack Frames

The next issue is the stack frame. Pascal supports nested subroutine definitions. A child subroutine (declared inside the parent subroutine) must be able to access all variables that are local to the parent in addition to its own local variables. In general, all variables declared at more outer scoping levels must be accessible to the inner subroutine. The subroutines in Figure B.1 illustrate the problems. `clio` can be accessed from inside all three subroutines. Other local variables can be accessed only from within the subroutines where they are declared, however. Similarly, `calliope` can't access `terpsichore` or `melpomene`. The situation is complicated, somewhat, by `erato` containing a recursive call to itself.

The stack, when the recursive call to `erato` is active, is shown in Figure B.2. There is one major difference between these stack frames and C stack frames: the introduction of a second pointer called the *static link*. The *dynamic link* is the old frame pointer, just as in C. The static link points, not at the previously active subroutine, but at the parent subroutine in the nesting sequence—in the declaration. Since `erato` and `thalia` are both nested inside `calliope`, their static links point at `calliope`'s stack frame. You can chase down the static links to access the local variables in the outer routine's

Static and dynamic links.

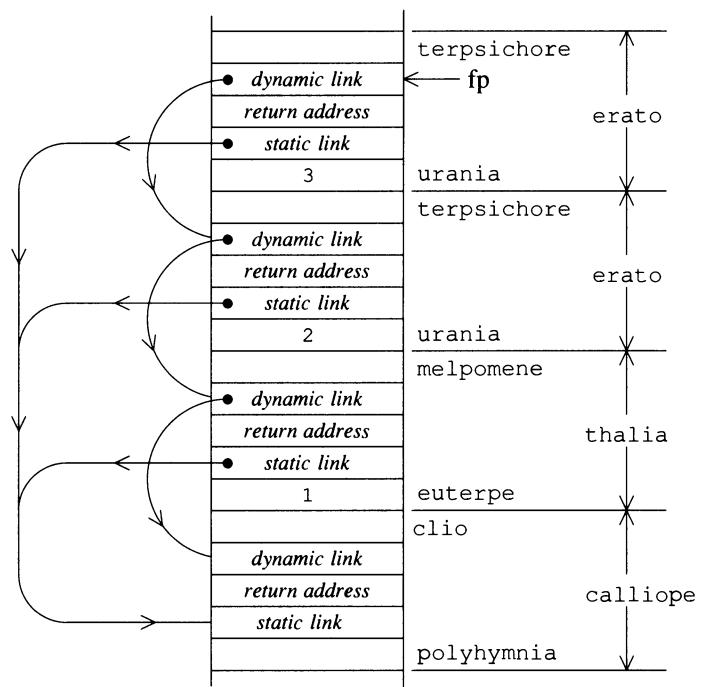
Figure B.1. Pascal Nested Subroutines

```

procedure calliope( polyhymnia: integer );
var
  clio: integer;
begin
  procedure erato( urania: integer );
  var
    terpsichore: integer;
  begin
    ...
    erato(3);
  end

  procedure thalia( euterpe: integer );
  var
    melpomene: integer;
  begin
    ...
    erato(2);
  end
  ...
  thalia(1);
end

```

Figure B.2. Pascal Stack Frames

stack frame. For example, `clio` can be accessed from `erato` with the following C-code:

```
r0.pp = WP(fp + 4);      /* r0 = static link */
_x    = W(r0.pp - 8);    /* x = clio           */
```

You can access `polyhymnia` from `erato` with:

```
r0.pp = WP(fp + 4);      /* r0 = static link */
_x    = W(r0.pp + 8);    /* x = clio           */
```

Though it's not shown this way in the current example, it's convenient for the frame pointer to point at the static, rather than the dynamic link to make the foregoing indirection a little easier to do. The static links can be set up as follows: Assign to each subroutine a *declaration level*, equivalent to the nesting level at which the subroutine is declared. Here, `calliope` is a level 0 subroutine, `erato` is a level 1 subroutine, and so forth. Then:

- If a subroutine calls a subroutine at the same level, the static link of the called subroutine is identical to the static link of the calling subroutine.
- If a subroutine at level N calls a subroutine at level $N+1$, the static link of the called subroutine points at the static link of the calling subroutine.
- If a subroutine calls a subroutine at a lower (more outer) level, use the following algorithm:

```
i = the difference in levels between the two subroutines;
p = the static link in the calling subroutine's stack frame;
while ( --i >= 0 )
    p = *p;
the static link of the called subroutine = p;
```

Note that the difference in levels (i) can be figured at compile time, but you must chase down the static links at run time. Since the static link must be initialized by the calling subroutine (the called subroutine doesn't know who called it), it is placed beneath the return address in the stack frame.

This organization can cause problems with `goto` statements. It's possible for a `goto` statement in `erato` to reference a label in `calliope`, for example. You can't just jump to the label in this case because the stack will be left with inactive stack frames on it, so you must treat the `goto` something like a `return`, adjusting the frame and stack pointers before making the jump—you must simulate return statements until the stack frame for the subroutine that contains the target label is active.

In order to make this simulation easier, Pascal delegates the responsibility of cleaning up the stack somewhat differently than C: it's the responsibility of the called function to delete the entire stack frame—both local variables and arguments—before returning. As an added advantage, the Pascal calling conventions tend to make somewhat smaller programs because the compiler doesn't have to generate an add-constant-to-stack-pointer directive after every subroutine call—programs usually have more subroutine calls than definitions. Finally, as a matter of convention rather than necessity, arguments to Pascal subroutines are pushed by the calling function in forward rather than reverse order. The leftmost argument is pushed first. This convention makes it difficult to write a Pascal function that takes a variable number of arguments unless the compiler provides some run-time mechanism for finding and identifying the type of the leftmost argument.

The Pascal `goto`.

Called subroutine cleans up stack. Less code is generated.

Arguments pushed in forward order.

A Grammar for C

This appendix summarizes the grammar used in Chapter Six. The organization here is different from that in Chapter Six in that it's more hierarchical. Line numbers in the following listing don't match the ones in Chapter 6. End-of-production actions are not shown, though imbedded actions are indicated with a {}. I've replaced the semicolon required at the end of a **yacc** or **occs** production with a blank line. The following abbreviations are used in the nonterminal names:

abs	abstract
arg(s)	argument(s)
const	constant
decl	declarator
def	definition
expr	expression
ext	external
opt	optional
param	parameter
struct	structure

Listing C.1. A Summary of the C Grammar in Chapter Six.

```

1 %union {                                     /* The value stack. */
2   char      *p_char;
3   symbol     *p_sym;
4   link       *p_link;
5   structdef  *p_sdef;
6   specifier   *p_spec;
7   value       *p_val;
8   int        num;           /* Make short if sizeof(int) > sizeof(int*). */
9   int        ascii;
10 }
11
12 %term  STRING      /* String constant. */          */
13 %term  ICON         /* Integer or long constant including '\t', etc. */
14 %term  FCON         /* Floating-point constant. */      */

```



Listing C.1. continued...

```

15 %term    TYPE          /* int char long float double signed unsigned short */
16           /* const volatile void */
17 %term    <ascii> STRUCT /* struct union */
18           ENUM        /* enum */
19
20 %term    <p_sym> TTYPE      /* type created with typedef */
21 %nonassoc <ascii> CLASS     /* auto extern register static typedef */
22 %nonassoc <p_sym> NAME
23 %nonassoc     ELSE
24
25 %term    RETURN GOTO
26 %term    IF ELSE
27 %term    SWITCH CASE DEFAULT
28 %term    BREAK CONTINUE
29 %term    WHILE DO FOR
30 %term    LC RC          /* { } */
31 %term    SEMI          /* ; */
32 %term    ELLIPSIS       /* ... */
33
34 %left   COMMA          /* , */
35 %right  EQUAL <ascii> ASSIGNOP /* = /= %= += -= <<= >>= &= |= ^= */
36 %right  QUEST COLON    /* ? : */
37 %left   OROR           /* || */
38 %left   ANDAND          /* && */
39 %left   OR              /* | */
40 %left   XOR             /* ^ */
41 %left   AND             /* & */
42 %left   <ascii> EQUOP    /* == != */
43 %left   <ascii> RELOP    /* <= >= < > */
44 %left   <ascii> SHIFTOP   /* >> << */
45 %left   PLUS MINUS     /* + - */
46 %left   STAR <ascii> DIVOP   /* * / % */
47 %right  SIZEOF <ascii> UNOP INCOP /* sizeof ! ~ ++ -- */
48 %left   LB RB LP RP <ascii> STRUCTOP /* [ ] ( ) . -> */

```

High-Level Definitions

```

49 program : ext_def_list { clean_up(); }
50
51 ext_def_list
52   : ext_def_list ext_def
53   | /* epsilon */
54
55 ext_def : opt_specifiers ext_decl_list {} SEMI
56   | opt_specifiers {} SEMI
57   | opt_specifiers funct_decl {} def_list {} compound_stmt
58
59 ext_decl_list
60   : ext_decl
61   | ext_decl_list COMMA ext_decl
62
63 ext_decl: var_decl
64   | var_decl EQUAL initializer
65   | funct_decl

```

Specifiers

```

66 opt_specifiers
67   : CLASS TTYPE
68   | TTYPE
69   | specifiers
70   | /* empty */

```



Listing C.1. continued...

```

71 specifiers
72     : type_or_class
73     | specifiers type_or_class
74
75 type    : type_specifier
76     | type type_specifier
77
78 type_or_class
79     : type_specifier
80     | CLASS
81
82 type_specifier
83     : TYPE
84     | enum_specifier
85     | struct_specifier

```

Enumerated Constants

```

87 enum_specifier
88     : enum name opt_enum_list
89     | enum LC enumerator_list RC
90
91 opt_enum_list
92     : LC enumerator_list RC
93     | /* empty */
94
95 enum     : ENUM
96
97 enumerator_list
98     : enumerator
99     | enumerator_list COMMA enumerator
100
101 enumerator
102     : name
103     | name EQUAL const_expr

```

Variable Declarators

```

104 var_decl: new_name           %prec COMMA
105     | var_decl LP RP
106     | var_decl LP var_list RP
107     | var_decl LB RB
108     | var_decl LB const_expr RB
109     | STAR var_decl          %prec UNOP
110     | LP var_decl RP
111
112 new_name: NAME
113 name   : NAME

```

Function Declarators

```

114 funct_decl
115     : STAR funct_decl
116     | funct_decl LB RB
117     | funct_decl LB const_expr RB
118     | LP funct_decl RP
119     | funct_decl LP RP
120     | new_name LP RP
121     | new_name LP {} name_list {} RP
122     | new_name LP {} var_list {} RP
123

```



Listing C.1. continued...

```

124 name_list
125     : new_name
126     | name_list COMMA new_name
127
128 var_list: param_declaration
129     | var_list COMMA param_declaration
130
131 param_declaration
132     : type var_decl
133     | abstract_decl
134     | ELLIPSIS

```

Abstract Declarators

```

135 abstract_decl
136     : type abs_decl
137     | TTYPE abs_decl
138
139 abs_decl: /* epsilon */
140     | LP abs_decl RP LP RP
141     | STAR abs_decl
142     | abs_decl LB RB
143     | abs_decl LB const_expr RB
144     | LP abs_decl RP
145

```

Structures

```

146 structSpecifier
147     : STRUCT opt_tag LC def_list RC
148     | STRUCT tag
149
150 opt_tag : tag
151     | /* empty */
152
153 tag     : NAME

```

Local Variables and Function Arguments

```

154 def_list: def_list def
155     | /* epsilon */
156
157 def     : specifiers decl_list {} SEMI
158     | specifiers SEMI
159
160 decl_list
161     : decl
162     | decl_list COMMA decl
163
164 decl    : funct_decl
165     | var_decl
166     | var_decl EQUAL initializer
167     | var_decl COLON const_expr          %prec COMMA
168     | COLON const_expr                  %prec COMMA
169
170 initializer
171     : expr                         %prec COMMA
172     | LC initializer_list RC
173
174 initializer_list
175     : initializer
176     | initializer_list COMMA initializer
177

```



Listing C.1. continued...

Statements

```
178 compound_stmt
179     : LC {} local_defs stmt_list RC {}
180
181 local_defs
182     : def_list
183
184 stmt_list
185     : stmt_list statement
186     | /* epsilon */
187
188 statement
189     : expr SEMI
190     | compound_stmt
191     | RETURN      SEMI
192     | RETURN expr  SEMI
193     | SEMI
194     | GOTO target SEMI
195     | target COLON {} statement
196     | SWITCH LP expr RP {} compound_stmt
197     | CASE const_expr COLON
198     | DEFAULT COLON
199     | IF LP test RP statement
200     | IF LP test RP statement ELSE {} statement
201     | WHILE LP test RP {} statement
202     | FOR LP opt_expr SEMI test SEMI {} opt_expr RP {} statement
203     | DO {} statement WHILE {} LP test RP SEMI
204     | BREAK SEMI
205     | CONTINUE SEMI
206
207 test    : {} expr
208     | /* empty */
209
210 target  : NAME
```

Expressions

```
211 opt_expr: expr
212     | /* epsilon */
213
214 const_expr
215     : expr %prec COMMA
216
217 expr    : expr COMMA {} non_comma_expr
218     | non_comma_expr
219
220 non_comma_expr
221     : non_comma_expr QUEST {} non_comma_expr COLON {} non_comma_expr
222     | non_comma_expr ASSIGNOP non_comma_expr
223     | non_comma_expr EQUAL    non_comma_expr
224     | or_expr
225
226 or_expr : or_list
227 or_list : or_list OROR {} and_expr
228     | and_expr
229
230 and_expr: and_list
231 and_list: and_list ANDAND {} binary
232     | binary
233
```

Listing C.1. continued...

```
234 binary : binary RELOP    binary
235     | binary EQUOP    binary
236     | binary STAR     binary
237     | binary DIVOP    binary
238     | binary SHIFTOP   binary
239     | binary AND      binary
240     | binary XOR      binary
241     | binary OR       binary
242     | binary PLUS     binary
243     | binary MINUS    binary
244     | unary
245
246 unary  : LP expr RP
247     | FCON
248     | ICON
249     | NAME
250     | string_const %prec COMMA
251     | SIZEOF LP string_const RP          %prec SIZEOF
252     | SIZEOF LP expr RP                 %prec SIZEOF
253     | SIZEOF LP abstract_decl RP        %prec SIZEOF
254     | LP abstract_decl RP unary        %prec UNOP
255     | MINUS unary                   %prec UNOP
256     | UNOP unary
257     | unary INCOP
258     | INCOP unary
259     | AND unary                   %prec UNOP
260     | STAR unary                  %prec UNOP
261     | unary LB expr RB           %prec UNOP
262     | unary STRUCTOP NAME        %prec STRUCTOP
263     | unary LP args RP
264     | unary LP      RP
265
266 string_const
267     : STRING
268     | string_const STRING
269
270 args   : non_comma_expr      %prec COMMA
271     | non_comma_expr COMMA args
```

LEX

This appendix is a user's manual for LEX,¹ a program that translates an input file made up of intermixed regular expressions and C source code into a lexical analyzer. Appendix E contains a user's manual for **occ**s, a compiler-compiler modeled after the UNIX **yacc** utility that translates an augmented, attributed grammar into a bottom-up parser.

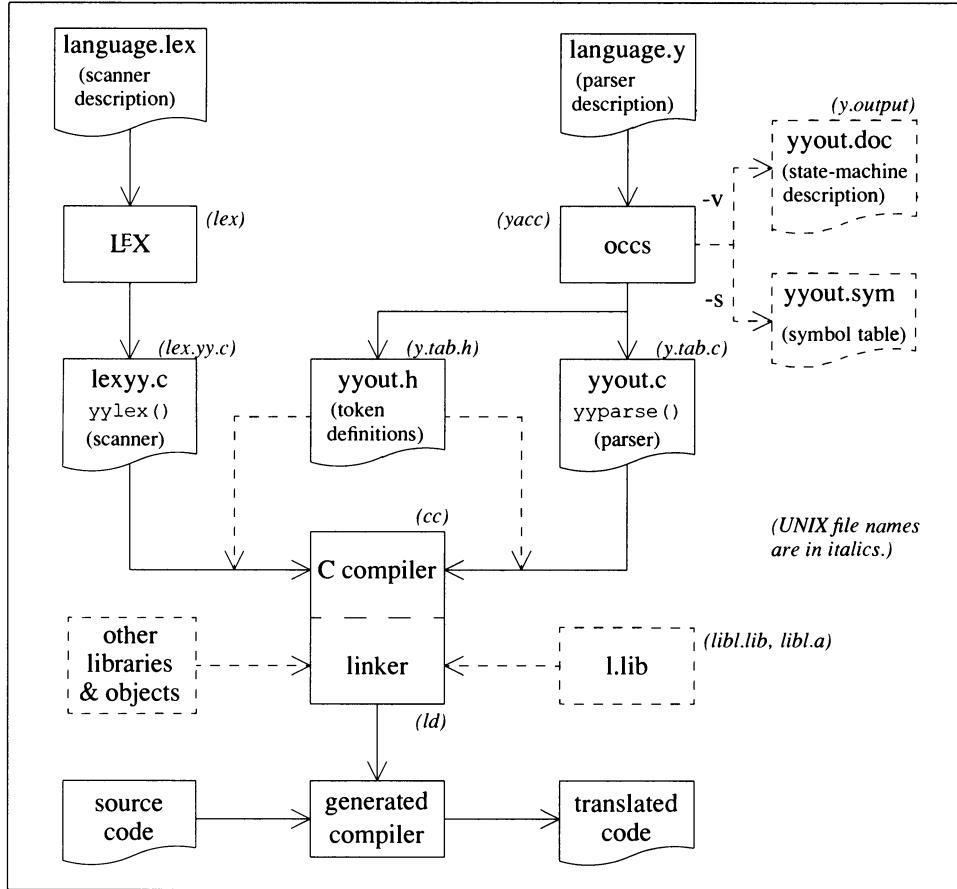
This (and the next) appendix describe LEX and **occ**s in depth. LEX and **occ**s are both very similar to the equivalent UNIX utilities, and you can use the UNIX **lex** and **yacc** rather than LEX and **occ**s to implement the compiler in Chapter Six, though you should read the UNIX documentation for these programs rather than this (and the next) appendix if you do. Footnotes are used to point out differences between LEX and **occ**s and their UNIX counterparts—you can ignore them if you're not interested in using the UNIX programs. In general, you should have little difficulty going from one system to the other.

D.1 Using LEX and Occs Together

LEX and **occ**s work together in an integrated fashion, though LEX can be used without **occ**s for many noncompiler applications. Both utilities can be viewed as C preprocessors. LEX takes as input files made up of interspersed LEX directives and C source code and outputs a table-driven lexical analyzer. **Occs** outputs a parser. The interaction between LEX and **occ**s is pictured in Figure D.1.

The LEX output file is called *lexyy.c*, and this file contains a tokenizer function called *yylex()*.² This function reads characters from the input stream, returning a token (and the associated lexeme) on every call. It is called by the **occ**s-generated parser every time a token is required. Similarly, **occ**s creates a file called *yyout.c*, which includes a parser subroutine called *yyparse()*.³ Somewhere in your program, there must be a call to

1. I'm calling it LEX to differentiate it from the UNIX program of the same name, which will be called **lex** throughout this manual. You can pronounce it as "leck" if you want to, but I wouldn't advise it.
2. UNIX **lex** creates a file called *lex.yy.c*; the lexical analyzer function is still called *yylex()*, however.
3. The **yacc** equivalent to *yyout.c* is *y.tab.c*; *yyparse()* is used by both programs.

Figure D.1. Creating a Compiler with L^EX and occs

`yparse()` to start up the parser.

The `l.lib` library contains various subroutines that are used by the lexical analyzer and should be linked into the final program.⁴ Note that `l.lib` includes a `main()` subroutine which loops, calling `yylex()` repetitively until `yylex()` returns zero. You can, of course, provide your own `main()`, and the one in the library will not be linked. [You must do this in the final compiler or else the parser will never be called. It's a common mistake to forget to provide a `main()`.]

There are several steps in creating a workable compiler using L^EX and occs. The first step is to write a lexical-analyzer specification as a L^EX input file, and to define a grammar in a way acceptable to occs (all this is discussed below and in Appendix E). You then run the input files through the appropriate programs, generating several `.c` and `.h` files.

L^EX creates `lexyy.c`, as was discussed earlier. Occs actually creates two source files: `yyout.c` contains source code for the parser and tables; `yyout.h` contains `#defines` for all the tokens.⁵

4. The UNIX programs use `/usr/lib/libl.lib` (alternately named `/usr/lib/libl.a` in some systems). Incorporate this library into your program with a `cc files... -l` directive when you link.

5. yacc calls the file `y.tab.h` and generates the file only if you specify `-d` on the command line. Occs always generates `yyout.h`, however.

One of these token values (and only these values) must be returned by your lexical analyzer for each input symbol.⁶ The special symbol `_EOI_`, always `#defined` as zero in `yyout.h`, is returned by the L^EX-generated analyzer at end of input. You should not return this value yourself as a token value, or `occs` (and `yacc`) will get confused.⁷ A typical development cycle with Microsoft C looks something like this:⁸

<code>vi compiler.lex</code>	<i>Create L^EX input specification.</i>
<code>vi compiler.y</code>	<i>Create occs input specification.</i>
<code>lex compiler.lex</code>	<i>Create lexyy.c from L^EX input specification.</i>
<code>occs -vD compiler.y</code>	<i>Create yyout.c and yyout.h from occs input specification.</i>
<code>cl -c lexyy.c</code>	<i>Compile everything and link.</i>
<code>cl -c yyout.c</code>	
<code>cl lexyy.obj yyout.obj -link l.lib</code>	

The various files generated by L^EX and `occs` are summarized in Table D.1. The `yyout.h` file is generally `#included` in the L^EX input specification, so you shouldn't compile `lexyy.c` until you've run `occs`.

D.2 The L^EX Input File: Organization

The remainder of this appendix describes the lexical-analyzer generator, L^EX. My implementation of L^EX is a proper subset of the similarly named UNIX utility.⁹ The input file is divided into three parts, separated from one another by `%%` directives:

```
definitions
%%
rules
%%
code
```

The `%%` must be the leftmost two characters on the line. The *definitions* section contains two things: macro definitions and user-supplied code. All code must be delimited

6. `yacc` also lets you use literal characters as tokens. If you put a `%term '+'` into the `yacc` input file, you can return an ASCII '+' from the lexical analyzer. `occs` doesn't support this mechanism. You'd have to put a `%term PLUS` in the `occs` input file and a `return PLUS;` statement in the lexical analyzer.

7. `yacc` doesn't define `_EOI_` for you—it nonetheless treats 0 as an end-of-input marker.

8. The UNIX development cycle looks like this:

```
vi compiler.lex
vi compiler.y
lex compiler.lex
yacc -vd compiler.y
cc -c lex.yy.c
cc -c y.tab.c
cc -o compiler y.tab.o lex.yy.o -ll
```

9. For those of you who are familiar with the UNIX program, the following features are not supported by L^EX:

- The / lookahead operator is not supported.
- The %START directive and the <name> context mechanism are not supported.
- The {N,M} operator (repeat the previous regular expression N to M times) is not supported.
- The REJECT mechanism is not supported (but `yymore()` is).
- The %S, %T, %C, and %R directives are not supported. Similarly, the internal-array-size modification directives %p, %n, %e, %a, %k, and %c) aren't supported.
- The internal workings of the output lexical analyzer (especially the input mechanism) bear virtually no relation to the UNIX program. You won't be able to hack the L^EX output in the same ways that you're used to.

Table D.1. L^EX and O^CCs Output Files

Generated by	O ^C Cs/L ^E X name	UNIX name	Command-line Switches	Description
L ^E X/lex	lexyy.c	lex.yy.c	none needed	C source code for lexical analyzer function, <code>yylex()</code> .
O ^C Cs/yacc	yyout.c	y.tab.c	none needed	C Source code for parser subroutine, <code>yyparse()</code>
	yyout.h	y.tab.h	occ s: none needed yacc : -d	# defines for the tokens.
	yyout.sym	not available	-s -S -D	Symbol-table showing various internal values used by the parser, including production numbers, TOKEN values, and so forth. This table is useful primarily for debugging.
	yyout.doc	y.output	occ s: -v -V yacc : -v	A human-readable description of the LALR state machine used by the parser. The file also contains copies of all the error and warning messages that are sent to the screen as the input is processed. It's useful for debugging the grammar. The <code>-V</code> version of the table holds LALR(1) lookaheads for all the items, and FIRST sets for all the nonterminals. The <code>-v</code> version contains lookaheads only for those items that trigger reductions, and does not contain FIRST sets.

with %{ and %} directives like this:

```
%{
#include <stdio.h>
extern int something; /* This code is passed to the output */
%}
%%
```

The percent sign must be the leftmost character on the line.¹⁰ Code blocks in the definitions section should contain **#includes**, external definitions, global-variable definitions, and so forth. In general, everything that would go at the top of a C program should also go at the top of the L^EX input file.

Anything outside of a %{ %} block is taken to be a macro definition. These consist of a name followed by white space (blanks or tabs), followed in turn by the macro's contents. For example:

10. UNIX lex supports directives like this:

```
%{ #include <stdio.h> %}
```

L^EX requires you to use the following, however:

```
%{ #include <stdio.h> %}
```

```
digit      [0-9]          /* decimal digit      */
hexdigit   [0-9a-fA-f]    /* hexadecimal digit  */
alpha      [a-zA-Z]        /* alphabetic character */
```

Comments.

Macro expansion is discussed further, below. C-style comments are permitted anywhere in the definitions section—they are ignored along with any preceding white space. C comments in a `%{ . . . %}` block are passed to the output, however.

Skipping the middle, *rules*, section for a moment, the third, *code* section is just passed to the output file verbatim. Generally, the initial definitions section should contain only code that is used in the rules section; moreover, it should contain only variable definitions, **extern** statements, and the like. All other code, including the subroutines that are called from the rules section, should be defined in the last, code section.

D.3 The L^EX Rules Section

Whitespace.

The *rules* in the middle part of the L^EX input specification have two components: a regular expression like the ones described both in Chapter Two and below, and C-source-code fragments that are executed when input that matches the regular expression is found. The regular expressions must start in the leftmost column, and they are separated from the code by white space (spaces or tabs). This means that you can't put a literal space character or tab into a regular expression. Use the `\s` or `\t` escape sequence for this purpose. You can also use double quotes—a " " matches a space—but the escape sequences are preferable because you can see them. The code part of the rule may take up several lines, provided that the lines begin with white space. That is, all lines that start with white space are assumed to be part of the code associated with the previous regular expression. A possible rule section looks like this:

```
%%
lama    printf("Priest\n");    /* A one-ell lama      */
llama   printf("Beast \n");    /* A two-ell llama   */
lll+ama printf("Fire! \n");    /* A three-ell lllama, four */
                                /* alarmer, and so forth. */
```

which, when given the input:

```
lama
llama
lllama
```

recognizes the various kinds of lamas.¹¹

Metacharacters.

D.3.1 L^EX Regular Expressions

The L^EX regular-expression syntax is essentially the one described in Chapter Two. It is summarized below, along with discussions of L^EX's peculiarities. Regular expressions are formed of a combination of normal characters and “metacharacters”, which have special meanings. The following characters are used as metacharacters—they are described in depth below:

```
* + ? { } [ ] ( ) . ^ $ "
```

11. See [Nash].

Note that the term “newline” is used in a nonstandard way in the following definitions. A *newline* here, is either an ASCII carriage return or a linefeed character (' \r' or ' \n'). This usage is consistent with most MS-DOS applications because the input characters are read in binary, not translated mode, so all input lines are terminated with a \r\n pair. By recognizing either '\n' or '\r' as a newline, **LEX**-generated code work can work in the UNIX environment, in MS-DOS translated, and in MS-DOS binary modes as well. If you want strict UNIX compatibility, you can use the *-u* command line switch discussed below to redefine newline to be a single linefeed character.

Newline, nonstandard definition.

- c A single character that is not used as a metacharacter forms a regular expression. The multiple-character sequences shown in Table D.2 also match single characters. These are known as *escape sequences*. Note that in MS-DOS, binary-mode input is used by **LEX**. This means that all lines are terminated with the two-character sequence \r\n. A \n matches the linefeed, but not the carriage return. It can be risky to test for an explicit \r or \n in the MS-DOS environment (but see ^ and \$, below). Double quotes can also be used to take away the meaning of metacharacters: "*?" matches an asterisk followed by a question mark.

Matching normal characters, escape sequences.

Table D.2. LEX Escape Sequences

Escape sequence	Matches	Recognized by UNIX
\b	backspace	yes
\f	formfeed	yes
\n	linefeed (end of line character)	yes
\r	carriage return	yes
\s	space	yes
\t	tab	yes
\e	ASCII ESC character ('\033')	no
\ddd	number formed of three the octal digits <i>ddd</i>	yes
\xdd	A number formed of the two hex digits <i>dd</i> . Two digits must be present—use \x02 for small numbers.	no
\^C	(backslash, followed by an up arrow, followed by a letter) matches a control character. This example matches Ctrl-C.	no
\c	\ followed by anything else matches that character; so, \. matches a dot, * matches an asterisk, and so forth.	yes

cc	Two regular expressions concatenated form a regular expression.	Concatenation.
e e	Two regular expressions separated by a vertical bar recognize either the expression to the left of the bar or the expression to the right of the bar.	OR operator.
(e)	Parentheses are used for grouping.	Grouping.
^	An up arrow anchors the pattern to beginning of line. If the first character in the expression is a ^, then the expression is recognized only if it is at the far left of a line.	Beginning-of-line anchor.
\$	A dollar sign anchors the pattern to end of line. If the last character in the expression is a \$, then the expression is recognized only if it is at the far right of a line.	End-of-line anchor.
.	A period (pronounced dot) matches any character except the newline (' \r' or ' \n' normally, ' \n' if <i>-u</i> is specified).	Match any character.

Character classes.

[...]

Brackets match any of the characters enclosed in the brackets. If the first character following the bracket is an up arrow (^), any character except the ones specified is matched. Only six characters have special meaning inside a character class:

{	Start of macro name
}	End of macro name
]	End of character class
-	Range of characters
^	Indicates negative character class
"	Takes away special meaning of characters up to next quote mark
\	Takes away special meaning of next character

Use \], \-, \\, and so forth, to put these into a class. Since other metacharacters such as *, ?, and + are not special here, the expression [*?+] matches a star, question mark, or plus sign. Also, a negative character class doesn't match a newline character—[^a-z] actually matches anything except a lower-case character or a newline. Use ((^a-z)|[\r\n]) to match the newline, too. Note that a negative character class must match a character—[^a-z]\$ does not match an empty line. The line must have at least one character on it, but that character may not be a lower-case letter.

Empty character classes.

L^EX supports two special character classes not supported by the UNIX lex. An empty character class ([]) matches all white space, with white space defined liberally as any character whose ASCII value is less than or equal to a space (all control characters and a space). An empty negative character class ([^]) matches anything but white space.

Closure operators.

* + ?

A regular expression followed by a * (pronounced star) matches that expression repeated zero or more times; a + matches one or more repetitions; a ? matches zero or one repetitions.

Greedy algorithm.

Note that L^EX uses a *greedy* algorithm to match a pattern. It matches the longest possible string that can satisfy the regular expression. Consequently, something like .* is dangerous (it absorbs all the characters on the line). The expression (.|\n)* tries to absorb an entire input file, but will probably cause an input-buffer overflow.

Macro expansion.

{name}

Braces are used to expand a macro name. (The {m,n} notation described in Chapter 2 is not supported by L^EX.) The *name* must be part of a macro definition in the first of the input file.

You can use the macros as an element of a regular expression by surrounding the name with curly braces. Given the earlier macro definitions:

```
digit      [0-9]          /* decimal digit      */
hexdigit   [0-9a-fA-f]    /* hexadecimal digit  */
alpha      [a-zA-Z]        /* alphabetic character */
%%
```

a hex number in C can be recognized with 0x{hexdigit}+. An alphanumeric character can be recognized with the following expression:

```
((alpha}|{digit}))
```

Macros nest, so you could do the following.

```
digit [0-9]
alpha [a-zA-Z]
alnum ((alpha}|{digit}))
```

You could do the same with:

```

digit 0-9
alpha a-zA-Z
alnum [{digit}{alpha}]

```

The ^ and \$ metacharacters work properly in all MS-DOS input modes, regardless of whether lines end with `\r\n` or a single `\n`. Note that the newline is not part of the lexeme, even though it must be present for the associated expression to be recognized. Use `\r\n` to put the end of line characters into the lexeme. (The `\r` is not required in UNIX applications, in fact it's an error under UNIX.) Note that, unlike the vi editor `^$` does not match a blank line. You'll have to use an explicit search such as `\r\n\r\n` to find empty lines.

The operator precedence is summarized in the Table D.3. All operators associate left to right.

The ^ and \$ metacharacters.

Operator Precedence.

Table D.3. Regular-Expression Operator Precedence

operator	description	level
()	parentheses for grouping	1 (highest)
[]	character classes	2
* + ?	closure: 0 or more, 1 or more, 0 or 1	3
cc	concatenation	4
	OR	5
^ \$	anchors to beginning and end of line	6 (lowest)

By default, input that doesn't match a regular expression is silently absorbed. You can cause an error message to be printed when bad input is encountered by putting the following into the definitions section:¹²

Input that doesn't match a regular expression.

```
%{
#define YYBADINP
%}
```

In general, it's a good idea to put one rule of the form

at the very end of the rules section. This default rule matches all characters not recognized by a previous expression, and does nothing as its action.

When two conflicting regular expressions are found in the input file, the one that comes first is recognized. The expression that matches the longest input string always takes precedence, however, regardless of the ordering of expressions in the file. Consider the following input file:

Conflicting regular expressions.

```
a.    return DOO;
aa+   return WHA;
```

Given the input aa, DOO will be returned because a. comes first in the file. Given the input aaa, the second expression is active because the longer string is matched.

12. UNIX lex behaves differently. It prints the unrecognized string to standard output and doesn't support YYBADINP.

D.3.2 The Code Part of the Rule

The code portion of the rule can be any legal C-code. It ends up in the output as part of a **switch** statement, so be careful with **breaks**. `yylex()` does not return unless an explicit **return** statement is included in one of the actions associated with a regular expression. On the other hand, a legal C statement must be attached to every regular expression, even if that statement is only a single semicolon.

Using local variables in actions.

The action code does not have to be surrounded with curly braces, but if you do so, you can use local variables in the code, as in:

```
expression {  
    int i;  
    ...  
}
```

The scope of the local variable is limited by the curly braces, as in normal C. Global variables must be defined in a `%{...%}` block in the definitions section.

If a vertical bar is used in place of an action, the action associated with the next rule is used. For example:

```
0x[0-9a-fA-F]+ |  
0[0-7]+ |  
[1-9][0-9]* return( NUMERIC_CONSTANT );
```

causes the **return** statement to be executed if any of the three expressions are matched.

Subroutines for use in actions.

Several subroutines, variables, and macros are available for use inside any of your actions (in either the rules or code section, but not in a `%{...%}` block in the definitions section). The subroutines are all in *l.lib*, which must be linked to your program. You can provide your own versions of all of them, however, in which case the ones in the library are not linked. The macros definitions are all imbedded in the LEX output file itself, and their definitions are discussed in depth in Chapter Two. LEX supports the following subroutines and macros:

Default `main()`.

```
void main()
```

The short `main()` subroutine shown in Listing D.1 is provided in *l.lib* as a convenience in testing lex-generated recognizers. It just loops, calling `yylex()` until `yylex()` returns zero. It's linked automatically if you don't provide a `main()` of your own. This default `main()` can cause problems if you forget to put a `main()` in an `occs`-generated compiler. In this case, `main()` is fetched from the library, and the parser is never activated.

Current lexeme.

```
char *yytext
```

This variable points at the lexeme that matches the current regular expression. The string is '\0' terminated.

Lexeme length.

```
int yylen
```

The number of characters in `yytext`, excluding the final '\0'.

Input line number.

```
int yylineno
```

The current input line number. If the lexeme spans multiple lines, it is the number of the last line of the lexeme.

Input a character to the current lexeme.

```
int input()
```

Read (and return) the next input character—the one following the last character in the lexeme. The new character is added to the end of the lexeme (`yytext`) by

Listing D.1. *yymain.c*—The main() Subroutine in *l.lib*

```

1 #include <stdlib.h>
2 #include <tools/debug.h>
3 #include <tools/l.h>
4
5 void main( argc, argv )
6 char **argv;
7 int argc;
8 {
9     /* A default main module to test the lexical analyzer.
10    */
11
12     if( argc == 2 )
13         ii_newfile( argv[1] );
14     while( yylex() )
15         ;
16     exit( 0 );
17 }
```

L^EX and **yyleng** is adjusted accordingly. Zero is returned at end of file, -1 if the lexeme is too long.¹³

ECHO

Echo current lexeme to the screen.

Print the current lexeme to a stream named **yyout**, which is initialized to **stdout**. This is just an alias for **fprintf(yyout, "%s", yytext)**. Do not use this macro if you intend to use the **occ**s interactive debugging environment described in Appendix E—it messes up the windows.

void output(int c)

Output a character.

Print **c** to the output stream. The **FILE** pointer used for this stream is called **yyout** and is initialized to **stdout**. You can change the file with **fopen()** if you like. Do not use this function if you intend to use the **occ**s interactive-debugging environment described in Appendix E—it messes up the windows.

void unput(int c)

Push back one character.

Pushes the indicated character back into the input stream. Note that **c** effectively overwrites **yytext[yyleng-1]** and then makes the lexeme one character smaller. Both **yyleng** and **yytext** are modified; **c** will be read as the next input character.¹⁴

YYERROR(char *s)

Print internal error message.

This macro is not supported by the **UNIX lex**. It is used to print internal error messages generated by the lexical analyzer itself. The default macro prints the string, **s**, to **stderr**. You will probably want to redefine this macro for **occ**s—put the following into the **L^EX** input file's definition section:

13. UNIX **lex** doesn't return -1 and it doesn't modify the **yytext** or **yyleng**; it just returns the next input character.

14. UNIX **lex** does not modify **yytext** or **yyleng**. It just pushes back the character.

	%{ #define YYERROR(str) yyerror("%s\n", str); }%
Push back multiple input characters.	void yyless(int n)
	Push <i>n</i> characters back into the input stream. Again, yytext and yylen are adjusted to reflect the pushed back characters. ¹⁵ You cannot push back more than yylen characters.

Initialize LEX parser.	void yy_init_lex()
	This subroutine is called by yylex() before it reads any input. You can use it to do any run-time initializations that are required before the lexical analyzer starts up. It's particularly handy if you want to initialize variables that are declared static in the definitions part of the LEX input file. You can define a yy_init_lex() that does the initializations at the bottom of the input file, and this function is called automatically. The default routine, in <i>l.lib</i> , does nothing. It's in Listing D.2. ¹⁶

Listing D.2. yyinitlx.c—Initialize LEX

```
1 void yy_init_lex( ) {} /* Default initialization routine--does nothing. */
```

Keep processing after expression is recognized.

Invoking this macro causes yylex() to continue as if the regular expression had not been recognized. Consider the following regular expression and action (which recognizes a subset of the legitimate string constants in C):

```
\\"[^\\"]*\\" { if( yytext[yylen-2] == '\\\' )
    yymore();
    else
        return STRING;
}
```

The regular expression recognizes a double quote ("") followed by anything except a second double quote repeated zero or more times. The quotes have to be escaped with a backslash because they're special characters to LEX. The problem here is escaped double quotes within strings, as in

```
"string with a \" in it"
```

The regular expression is satisfied when the first double quote is encountered. The action code looks at the character preceding the ending quote mark (yytext[yylen-2]). If that character is a backslash, yymore() is called, and more input will be processed as if the offending double quote had not been seen. The same code is executed a second time when the third double quote is found; but here, there's no preceding backslash so the **return** STRING is executed.

15. UNIX lex does not modify yytext or yylen; it just pushes back the characters.

16. UNIX lex does not support this mechanism. You're on your own when it comes to initializations.

int yywrap()

End-of-input processing.

This subroutine is called by the lexical analyzer at end of file. The name stands for “go ahead and wrap up the processing”, so `yywrap()` returns true if processing should terminate [in which case `yylex()` returns zero.] The default routine, in `l.lib`, just returns 1. It’s in Listing D.3. You can provide your own `yywrap()` however, and yours will be linked instead of the library routine. Typically, you’d use `yywrap()` to handle a series of input files. It would open the next input file in the series with each call, returning zero until there were no more files to open, whereupon it would return one. There’s an example in Listing D.4.

Listing D.3. `yywrap.c`— Library Version

```

1  int      yywrap()          /* yylex() halts if 1 is returned */
2  {
3      return( 1 );
4 }
```

Listing D.4. A User-Defined `yywrap()`

```

1  int  Argc;                      /* Copy of argc as passed to main(). */
2  char **Argv;                  /* Copy of argv as passed to main(). */
3
4  yywrap()
5  {
6      if( --Argc >= 0 )
7      {
8          if( ii_newfile( *Argv ) != -1 )
9          {
10             ++Argv;
11             return 0;           /* New file opened successfully. */
12         }
13         fprintf(stderr, "Can't open %s\n", *Argv );
14     }
15     return 1;
16 }
17
18 main( argc, argv )
19 int    argc;
20 char   **argv;
21 {
22     Argc = argc - 1;
23     Argv = argv + 1;
24     ii_newfile( *argv )
25     while( yylex() )
26         ;
27 }
```

YYBADINP

If this macro is **#defined** in a `%{...%}` block in the definitions section, an error message is printed when an input sequence that’s not recognized by any regular expression is encountered at run time. Otherwise, the unrecognized input is silently ignored at run time.¹⁷

Control error message for unrecognized input character.

Low-level input functions.

Low-level input routine.

Low-level lookahead.

Open new input file.

Flush low-level input buffers.

Change system-level input functions.

The longest permissible lexeme is 1024 bytes.¹⁸ An error message is printed if you try to read more characters. If a lexeme is likely to exceed this length, you'll have to work with one of the low-level input functions that L^EX itself uses, synopsized below, and described in depth in Chapter Two. None of these functions are supported by UNIX lex.

```
int ii_input()
```

This is a somewhat more flexible input function than input(). Normally, it returns the next input character (and stretches the lexeme one character). It returns zero on end of file and -1 if it can't read a character because the input buffer is full [see ii_flushbuf()]. Note that yylen is not modified by ii_input().

```
int ii_lookahead(int n)
```

You can use this routine to look ahead in the input without actually reading a character. ii_lookahead(0) returns the last character in the lexeme. ii_lookahead(1) returns the next character that will be read. ii_lookahead(-1) returns the penultimate character in the lexeme. The maximum forward lookahead is 32 characters, and you can look back only as far as the beginning of the current lexeme.

```
int ii_newfile(char *name)
```

Use this routine to open a new input file. The *name* is the file name. The routine returns a file descriptor for the open file or -1 if it can't open the file. As is the case with normal input functions, the global variable yyerror will hold an error code identifying the problem. This routine automatically closes the current input file before opening the new one, so it's actually an error if you close the current input file explicitly before calling ii_newfile().¹⁹

```
int ii_flushbuf()
```

This routine flushes the input buffer, discarding all characters currently in it and destroying the current lexeme. It can be used to continue reading when ii_input() returns -1. Note that yylex, yylen, and yylineno are all invalid after this routine is called. It's usually used only when you're doing something like absorbing a long comment—when the contents of the lexeme are immaterial. It returns 1 if everything's okay, 0 at end of file.

```
int ii_io(int (*open)(), int (*close)(), int (*read)())
```

You can use this routine to change the low-level, unbuffered I/O functions used by the input routines. It takes three function pointers as arguments, the first is a pointer to an open function, the second to a close function, and the third to a read function. These should work just like the standard UNIX I/O functions with the same name, at least in terms of the external interface. The open function is called

17. UNIX lex does not support this mechanism. It just writes the unrecognized characters to standard output.

18. 128 bytes is typical in UNIX lex.

19. UNIX lex does not use this input mechanism; rather, it uses the buffered I/O system for its input and gets input from a stream called yyin. This stream is initialized to stdin, but you can modify it using fopen(). You must use ii_input() to change the L^EX input stream, however.

as follows:

```
#include <fcntl.h>
#ifndef MSDOS
#define O_MODE (O_RDONLY|O_BINARY)
#else
#define O_MODE (O_RDONLY)
#endif
...
int fd;
...
if( (fd = (*open) (name, O_MODE)) != -1 )
{
    /* file open was successful */
}
return fd;
```

It should return an `int`-sized number that can't be confused with standard input (anything but zero), and that same number will be passed to the read and close functions as follows:

```
int got;
char *load_characters_here; /* base address of the buffer */
int read_this_many_bytes; /* number of characters requested */
int got; /* number of characters actually read */
int fd; /* value returned from previous open */

...
if((got=(*read)(fd, load_characters_here, read_this_many_bytes))==-1)
    process_end_of_file();
...
(*close)( fd );
```

All variables and subroutines used both by L^EX and its input functions start with one of the following character sequences:²⁰

yy Yy YY ii_

Don't start any of your own symbols with these characters.

Names start with yy Yy
YY ii_.

D.4 L^EX Command-Line Switches

L^EX takes various command-line switches that modify its behavior. These are summarized in Table D.4²¹

L^EX automatically compresses the tables that it generates so that they'll take up less room in the output. A ten-to-one compression rate is typical. The default compression method eliminates equivalent rows and columns from the tables and is best for most practical applications. The `-c` and `-f` switches can be used to control the table-compression, however. The `-f`(for fast) switch eliminates compression entirely, yielding much larger tables, but also giving a faster lexical analyzer. The `-c` switch changes the compression algorithm to one in which each character/next-state pair is stored as a two-byte object. This method typically (though not always) gives you smaller tables and a

Controlling table compression, `-c`, `-f`.

20. You have to worry only about the yy prefix with `lex`.

21. Of these, only `-t` and `-f` are supported by UNIX `lex`. The `lex -v` switch prints a one-line summary of internal statistics only—the L^EX `-v` is more verbose. The `lex -n` switch suppress the verbose-mode output, and is not supported by L^EX.

Table D.4. L^EX Command-Line Switches

Switch	Description
<code>-cN</code>	Use pair compression, N is the threshold. The default N is 4.
<code>-f</code>	For fast. Don't compress tables.
<code>-h</code>	Suppress header comment that describes state machine.
<code>-H</code>	Print the header only.
<code>-l</code>	Suppress <code>#line</code> directives in the output.
<code>-mname</code>	Use <i>name</i> as template-file name rather than <i>lex.par</i> .
<code>-t</code>	Send output to standard output instead of <i>lexyy.c</i>
<code>-u</code>	UNIX mode (. is everything but <code>\n</code>)
<code>-v</code>	Verbose mode, print various statistics.
<code>-V</code>	More verbose, print internal diagnostics as L ^E X runs.

Header comment, `-h`, `-H`.

Suppressing `#line` directives, `-l`.

Choose an alternate driver file, `-m`.

The LIB environment.

Use standard output, `-t`.
UNIX-compatible newline, `-u`.

Verbose-mode output, `-v`, `-V`.

slower lexical analyzer than the default method. The `-c` switch takes an optional numeric argument (`-c5`) that specifies the threshold above which the pair compression kicks in for a given row of the table. If the row has more than the indicated number of nonerror transitions, it is not compressed. The default threshold, if no number is specified, is four. In any event, the `-v` switch can be used to see the actual table size, so you can decide which method is most appropriate in a particular situation.

The `-h` and `-H` switches are used to control the presence of a larger header comment that is output at the top of *lexyy.c*. This comment describes the state machine that L^EX uses to recognize the regular expressions. A lower-case *h* suppresses this comment. An upper-case *H* suppresses all output except the comment.

L^EX automatically generates `#line` directives in the output file. These cause the compiler to print error messages that reference lines in the the original L^EX input file rather than the output file. They are quite useful when you're trying to track down syntax errors. The `#line` can confuse source-level debuggers, however, so `-l` (that's an ell) is provided to eliminate them.

L^EX itself creates only a small part of the output file. The majority is stored in a *template file* called *lex.par*, discussed in depth in Chapter 2. L^EX searches for this file, first in the current directory and then along a path specified by semicolon-delimited directory names in the LIB environment. The `-mname` switch can be used to specify an explicit file name to use as the template. There's no space between the *m* and the name.

The `-t` switch causes L^EX to send output to standard output rather than *lexyy.c*.

The `-u` switch changes the definition of a “newline” to be consistent with UNIX rather than MS-DOS. Normally a newline is either a carriage return or linefeed ('`\r`' or '`\n`'). In UNIX mode, however, a newline is a single linefeed. This definition affects the `.` operator, which matches anything except a newline, and a negative character class, which matches anything except a newline or one of the specified characters.

The `-v` switch causes L^EX to print various statistics, such as the output table sizes, to standard output. `-V` forces more-verbose output, which describes the internal workings of the program as it runs. This last switch is useful primarily for debugging L^EX itself; it's also handy if you want to see a running commentary of the subset construction used to create a DFA from an NFA, described in Chapter Two.

D.5 Limits and Bugs

Table-size limits.

Several limits are imposed on the L^EX input file. In addition, several internal limits will affect the workings of L^EX. These are summarized in Table D.5. There are two

limits on the action components to the rules. No single rule (including the action) can be longer than 2,048 bytes, and the space used by all the actions combined cannot exceed 20,480 bytes. If either of these limits are exceeded, fix the problem by moving some of the imbedded code into subroutines declared in the third part of the L^EX input specification.

Table D.5. L^EX Limits

Maximum space available for all actions combined	20,480	bytes
Maximum characters in a single rule (including the action)	2,048	bytes
Maximum number of NFA states	512	states
Maximum number of DFA states	254	states

The limits on NFA and DFA states refer to the maximum size of internal tables used by L^EX (see Chapter Two). There are always fewer DFA than NFA states. If you exceed either of these limits, you'll either have to reduce the number of regular expressions in the input file or simplify the expressions. Two simplification techniques can be used with little difficulty:

- (1) Character classes take up less room than the | operator so they should be used whenever possible; [012] is preferable to (0|1|2). The character class uses two states; the expression with | operator uses ten.
- (2) Use multiple rules with shared actions rather than the OR operator if possible. This rule:

```
abc      |
def      printf("abc or def");
```

requires one fewer NFA state than this one:

```
abc|def printf("abc or def");
```

- (3) One of the least efficient ways to use L^EX is to recognize long strings that do not contain metacharacters. A simple string uses roughly as many NFA states as it has characters—"0123456789" requires 11 states, but a closure operator or character class uses only a few states. These sorts of strings can cause problems when you try to do something like this:

```
<Ctrl-A>      printf("<Ctrl-A>\n");
<Ctrl-B>      printf("<Ctrl-B>\n");
<Ctrl-C>      printf("<Ctrl-C>\n");
<Ctrl-D>      printf("<Ctrl-D>\n");
<Ctrl-E>      printf("<Ctrl-E>\n");
<Ctrl-F>      printf("<Ctrl-F>\n");
<Ctrl-G>      printf("<Ctrl-G>\n");
<Ctrl-H>      printf("<Ctrl-H>\n");
<Ctrl-I>      printf("<Ctrl-I>\n");
<Ctrl-J>      printf("<Ctrl-J>\n");
<Ctrl-K>      printf("<Ctrl-K>\n");
<Ctrl-L>      printf("<Ctrl-L>\n");
```

Simplifying expressions
to get smaller internal
table size.

Use character classes,
not |.

Use shared actions,
not |.

Don't use long literal
matches.

Though the final state machine uses only 31 DFA states, 121 NFA states are required to produce this machine. A better solution is:

```
%%
<Ctrl-[A-L]> switch( yytext[6] )
{
    case 'A': printf("<Ctrl-A>\n"); break;
    case 'B': printf("<Ctrl-B>\n"); break;
    case 'C': printf("<Ctrl-C>\n"); break;
    case 'D': printf("<Ctrl-D>\n"); break;
    case 'E': printf("<Ctrl-E>\n"); break;
    case 'F': printf("<Ctrl-F>\n"); break;
    case 'G': printf("<Ctrl-G>\n"); break;
    case 'H': printf("<Ctrl-H>\n"); break;
    case 'I': printf("<Ctrl-I>\n"); break;
    case 'J': printf("<Ctrl-J>\n"); break;
    case 'K': printf("<Ctrl-K>\n"); break;
    case 'L': printf("<Ctrl-L>\n"); break;
}
%%
```

Use lookup tables to disambiguate lexemes.

Problems with anchors.

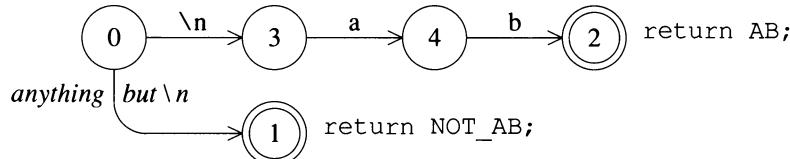
which requires only 11 NFA states.

Keyword recognition presents a similar problem that can be solved in much the same way. A single expression can be used to recognize several similar lexemes, which can then be differentiated with a table lookup or calls to `strcmp()`. For example, all keywords in most languages can be recognized by the single expression `[a-zA-Z_]+`. Thereafter, the keywords can be differentiated from one another with a table look up. There's an example of this process below.

A bug in LEX affects the way that the default . action is handled if the beginning-of-line anchor (^) is also used in the same input specification. The problem is that an explicit match of \n is added to the beginning of all patterns that are anchored to the start of line (in order to differentiate ^x from x\$ from ^x\$ from x—it should be possible to have all four patterns in the input file, each with a different accepting action). The difficulty here is that an input file like this:

```
%%
^ab  return AB;
.   return NOT_AB;
%%
```

generates a machine like this:



The problem lies in the way that the LEX-generated state machine works. When it fails out of a nonaccepting state, the machine backs up to the most-recently-seen accepting state. If there is no such state, the input lexeme is discarded. Looking at the previous machine, if you are in State 4, the machine fails if a character other than b is in the input. Not having seen an accepting state, the machine discards the newline and the a rather than performing the default `return NOT_AB` action. You can correct the problem by making State 3 an accepting state that also executes the `return NOT_AB` action. Do this by modifying the input file as follows:

```

%%  

^ab  return AB;  

\n  |  

.   return NOT_AB;  

%%

```

D.6 Example: A Lexical Analyzer for C

Listing D.5. shows a sample L^EX input file for a C lexical analyzer. Listing D.6, *yyout.h*, contains the token definitions. Various macros are defined on lines 11 to 17 of Listing D.5. The *suffix* macro is for recognizing characters that can be appended to the end of a numeric constant, such as 0x1234L. Note that all characters from an ASCII NUL up to a space (that is, all control characters) are recognized as white space.

Listing D.5. *c.lex*—L^EX Input File for a C Lexical Analyzer

```

1  %{
2  /* Lexical analyzer specification for C. Note that CONSTANTS are
3   * all positive in order to avoid confusions (to prevent a-l from
4   * being interpreted as NAME CONSTANT rather than NAME MINUS CONSTANT).
5   */
6
7  #include "yyout.h"
8  #include <search.h>      /* Function prototype for bsearch() */
9  %}
10
11 let     [_a-zA-Z]          /* Letter */                                */
12 alnum   [_a-zA-Z0-9]        /* Alphanumeric character */                         */
13 h       [0-9a-fA-F]        /* Hexadecimal digit */                                */
14 o       [0-7]               /* Octal digit */                                */
15 d       [0-9]               /* Decimal digit */                                */
16 suffix  [UuLl]             /* Suffix in integral numeric constant */           */
17 white   [\x00-\s]          /* White space */                                */
18
19 %%
20 "/*"                {
21     int i;
22
23     while( i = ii_input() ) {
24
25         if( i < 0 )
26             ii_flushbuf(); /* Discard Lexeme */
27
28         else if( i == '*' && ii_lookahead(1) == '/' )
29         {
30             ii_input();
31             break;          /* Recognized comment */
32         }
33     }
34
35     if( i == 0 )
36         yyerror( "End of file in comment\n" );
37     }
38
39 \"(\\".|[^\"])*\"    return STRING;
40 \"(\\".|[^\"])*[\r\n]  yyerror("Adding missing \" to string constant\n");
41                           return STRING;
42

```



Listing D.5. continued...

```

43  '.'          |      /* 'a', 'b', etc.          */
44  '\.'         |      /* '\t', '\f', etc.          */
45  '\\{o}({o}{o})?' |      /* '\123', '\12', '\1'        */
46  '\\x{h}({h}{h})?' |      /* '\x123', '\x12', '\x1'        */
47  0{o}*{suffix}? |      /* 0, 01, 012, 012L, etc.      */
48  0x{h}+{suffix}? |      /* 0x1, 0x12, 0x12L, etc.      */
49  [1-9]{d}*{suffix}? |      return ICON; /* 123, 123L, etc.          */
50
51 ({d}+|{d}+\.{d}*|{d}*\.{d}+) ([eE] [-+]?{d}+)?[fF]? |      return FCON ;
52
53 "("          |      return LP;
54 ")"          |      return RP;
55 "{"           |      return LC;
56 "}"           |      return RC;
57 "["           |      return LB;
58 "]"           |      return RB;
59
60 "->"        |
61 ". ."         |      return STRUCTOP;
62 "++"          |
63 "--"          |      return INCOP;
64 "[^!]"        |      return UNOP;
65
66 "*"
67 [ /%]
68 "+"
69 "-"
70 <<|>>
71 [<>]=?
72 [=!] =
73 [*/%+-&|^]=
74 (<<|>>)=
75 "="
76 "&"
77 "^"
78 "|"
79 "&&"
80 "| |"
81 "?"
82 ":"           |
83 ","
84 ";"
85 {let}{alnum}*
86 {white}+
87 .
88 %%
89
90 /*-----*/
91
92 typedef struct          /* Routines to recognize keywords */
93 {
94     char *name;
95     int val;
96 }
97 KWORD;
98
99 KWORD Ktab[] =          /* Alphabetic keywords */
100 {
101     { "auto",     CLASS      },
102     { "break",    BREAK      },

```



Listing D.5. continued...

```

103     { "case",      CASE      },
104     { "char",       TYPE      },
105     { "continue",   CONTINUE  },
106     { "default",    DEFAULT   },
107     { "do",         DO        },
108     { "double",    TYPE      },
109     { "else",       ELSE      },
110     { "extern",    CLASS     },
111     { "float",     TYPE      },
112     { "for",        FOR      },
113     { "goto",      GOTO     },
114     { "if",         IF        },
115     { "int",        TYPE      },
116     { "long",       TYPE      },
117     { "register",  CLASS     },
118     { "return",    RETURN   },
119     { "short",     TYPE      },
120     { "sizeof",    SIZEOF   },
121     { "static",    CLASS     },
122     { "struct",    STRUCT   },
123     { "switch",   SWITCH   },
124     { "typedef",  CLASS     },
125     { "union",     STRUCT   },
126     { "unsigned",  TYPE      },
127     { "void",      TYPE      },
128     { "while",     WHILE     }
129 };
130
131 int      cmp( a, b )
132 KWORD   *a, *b;
133 {
134     return strcmp( a->name, b->name );
135 }
136
137 int      id_or_keyword( lex ) /* Do a binary search for a */
138 char    *lex;                /* possible keyword in Ktab */
139 {                            /* Return the token if it's */
140     KWORD      *p;           /* in the table, NAME */
141     KWORD      dummy;        /* otherwise. */
142
143     dummy.name = lex;
144     p = bsearch( &dummy, Ktab, sizeof(Ktab)/sizeof(KWORD),
145                  sizeof(KWORD), cmp );
146     return( p ? p->val : NAME );
147 }
```

The rule on lines 20 to 37 of Listing D.5 handles a comment. The main problem here is the potential comment length. A single regular expression can't be used to recognize a comment because various internal limits on lexeme length will be exceeded by long comments. So, the code looks for the beginning-of-comment symbol, and then sits in a **while** loop sucking up characters until an end-of-comment is encountered. The *ii_flushbuf()* call on line 26 forces an internal buffer flush (discarding the partially-collected lexeme) if the input buffer overflows. Note that this comment processing isn't portable to UNIX **lex**. An alternate strategy is shown in Listing D.7. This second method works fine under UNIX, but disallows comments longer than 1024 bytes if used with my **LEX**.

Comments.

Listing D.6. *yyout.h*—C Token Definitions

line	/*	token	value	lexeme	*/
1					/*
2					
3	#define	_EOI	0	/* end-of-input symbol	*/
4	#define	NAME	1	/* identifier	*/
5	#define	STRING	2	/* "string constant"	*/
6	#define	ICON	3	/* integer constant or character constant	*/
7	#define	FCON	4	/* floating-point constant	*/
8	#define	PLUS	5	/* +	*/
9	#define	MINUS	6	/* -	*/
10	#define	STAR	7	/* *	*/
11	#define	AND	8	/* &	*/
12	#define	QUEST	9	/* ?	*/
13	#define	COLON	10	/* :	*/
14	#define	ANDAND	11	/* &&	*/
15	#define	OROR	12	/*	*/
16	#define	RELOP	13	/* > >= < <=	*/
17	#define	EQUOP	14	/* == !=	*/
18	#define	DIVOP	15	/* / %	*/
19	#define	OR	16	/*	*/
20	#define	XOR	17	/* ^	*/
21	#define	SHIFTOP	18	/* >> <<	*/
22	#define	INCOP	19	/* ++ --	*/
23	#define	UNOP	20	/* ! ~	*/
24	#define	STRUCTOP	21	/* . ->	*/
25	#define	TYPE	22	/* int, long, etc.	*/
26	#define	CLASS	23	/* extern, static, typedef, etc.	*/
27	#define	STRUCT	24	/* struct union	*/
28	#define	RETURN	25	/* return	*/
29	#define	GOTO	26	/* goto	*/
30	#define	IF	27	/* if	*/
31	#define	ELSE	28	/* else	*/
32	#define	SWITCH	29	/* switch	*/
33	#define	BREAK	30	/* break	*/
34	#define	CONTINUE	31	/* continue	*/
35	#define	WHILE	32	/* while	*/
36	#define	DO	33	/* do	*/
37	#define	FOR	34	/* for	*/
38	#define	DEFAULT	35	/* default	*/
39	#define	CASE	36	/* case	*/
40	#define	SIZEOF	37	/* sizeof	*/
41	#define	LP	38	/* ((left parenthesis)	*/
42	#define	RP	39	/*)	*/
43	#define	LC	40	/* { (left curly)	*/
44	#define	RC	41	/* }	*/
45	#define	LB	42	/* [(left bracket)	*/
46	#define	RB	43	/*]	*/
47	#define	COMMA	44	/* ,	*/
48	#define	SEMI	45	/* ;	*/
49	#define	EQUAL	46	/* =	*/
50	#define	ASSIGNOP	47	/* += -= etc.	*/

String constants.

The next two rules, on lines 39 to 41 handle string constants. The rules assume that a string constant cannot span a line, so the first rule handles legal constants, the second rule handles illegal constants that contain a newline rather than a terminating close quote. (ANSI is unclear about whether newlines are permitted in string constants. I'm assuming that they are not permitted because string concatenation makes them unnecessary. Consequently, a hard error is printed here rather than a warning.)

Listing D.7. A UNIX Comment-Processing Mechanism

```

1  %%
2  "/*"          {
3      int i, lasti = 0;
4
5      while( (i = input()) && i != EOF )
6      {
7          if( i == '/' & last_i == '*' );
8              break;
9      }
10
11     if( i == 0 || i == EOF )
12         yyerror( "End of file in comment\n" );
13
14 %%
```

Most of the complexity comes from handling backslashes in the string correctly. The strategy that I used earlier:

```

\"[^\\]*\"    if( yytext[yyleng-2] == '\\\\' )
                yymore();
else
    return STRING;
```

doesn't work in a real situation because this expression can't handle strings in which the last character is a literal backslash, like "\\\\". It also can't handle newlines in string constants, for which an error should be printed. The situation is rectified on line 39 of Listing D.5 by replacing the character class in the middle of \"[^\\]*\" with the following subexpression (I've added some spaces for clarity):

```
( \\\\. | [^\\"] )*
```

This subexpression recognizes two things: two-character sequences where the first character is a backslash, and single characters other than double quotes. The expression on line 40 handles the situation where the closing quote is missing. It prints an error message and returns a STRING, as if the close quote had been found.

All the rules on lines 43 to 49 recognize integer constants of some sort. The first five are for the various character constants and the last four are for numeric constants. Floating-point constants are recognized on line 51.

The rule on line 85 recognizes both identifiers and keywords. I've done this to minimize the table sizes used by the LEX-generated analyzer. The id_or_keyword() subroutine is used to distinguish the two. It does a binary search in the table declared on lines 99 to 129, returning the appropriate token if the current lexeme is in the table, otherwise it returns NAME. The bsearch() function on line 144 does the actual search. It is a standard ANSI function that's passed a key, a pointer to a table, the number of elements in the table, the size of one element (in bytes), and a pointer to a comparison function. This last function is called from bsearch(), being passed the original key and a pointer to one array element. It must compare these, but otherwise work like strcmp(). The comparison function used here is declared on lines 131 to 135. The remainder of the rules are straightforward.

Using lookup table to
differentiate identifiers
from keywords.

D.7 Exercises

- D.1. Newlines in string constants are not handled particularly well by the LEX input file presented in the previous section. The current file assumes that newlines are not permitted in constants, so it inserts a close quote when a newline is encountered. Rewrite the regular expressions for string constants so that they will do the following:
- Print a warning if a newline is found in a constant, and then take the next line as a continuation of the previous line. That is, the close quote should not be inserted by the lexical analyzer; rather, it should recognize multiple-line strings.
 - If the newline in the string constant is preceded by a backslash, however, the warning message is suppressed.
- D.2. The UNIX *vi* editor supports a *tags* feature. When the cursor is positioned on a subroutine name, typing a Ctrl-] causes the current file to be saved, the file that contains the indicated subroutine is then read into the editor, and the cursor is positioned at the first line of the subroutine. *Vi* accomplishes this feat with the aid of supplemental file called *tags*, which is created by a utility called *ctags*. The tags file consists of several lines, each having the following fields:

subroutine_name	file_name	search_pattern
-----------------	-----------	----------------

The three fields are separated by tabs, and the search pattern consists of the actual first line of the subroutine, surrounded by / and \$/. For example:

```
main      main.c  /`main( argc, argv )$/
```

Using LEX, write a version of *ctags* that recognizes all subroutine declarations and all subroutine-like macro definitions in a C program. It should not recognize external declarations. That is, none of the following should generate an entry in the output file:

```
extern      larry ();
char        *curly( int x );
PRIVATE double  moe   ( void );
```

You can assume that subroutine declarations will follow a standard formatting convention. Hand in both the program and documentation describing what the input must look like for a declaration to be recognized.

- D.3. Write a version of *ctags* for Pascal programs.
- D.4. Using LEX, write a version of the UNIX calendar utility. Your program should read in a file called *calendar* and print all lines from that file that contain today's date somewhere on the line. All the following forms of dates should be recognized:

7/4/1776	
7/4/76	Assume 20th century, so this entry is for 1976.
7/4	Assume the current year.
7-4-1776	
7-4-76	
7-4	
1776-7-4	Since 1776 can't be a month, assume European ordering.
July 4, 1776	

July 4 1776

July 4

Jul 4

Three-character abbreviations for all month names are supported.

Jul. 4 The abbreviation can be followed by an optional period.

4 July

4 July, 1776

4 July, 76

4 July, '76

4 Jul. '76

LLama and Occs

LLama and **occ**s are compiler-generation tools that translate an augmented, attributed grammar into a parser. They are essentially C preprocessors, taking as input a file containing intermixed C source code, grammatical rules, and compiler directives, and outputting the C source code for a parser subroutine.

This appendix is a user's manual for both **LLama** and **occ**s, the two compiler compilers developed in Chapters Four and Five. (**Occs** stands for "the Other Compiler-Compiler System". It's pronounced "ox.") An understanding of the theoretical parts of these chapters (sections marked with asterisks in the table of contents) will be helpful in reading the current Appendix. Similarly, a discussion of how to use **occ**s and **LLama** in conjunction with **LEX** is found at the beginning of Appendix D, which you should read before continuing.

Both **LLama** and **occ**s are modeled after the UNIX **yacc** utility, though **LLama** builds a top-down, LL(1) parser while **occ**s builds a bottom-up, LALR(1) parser. The compiler in Chapter Six can be constructed both with **occ**s or with **yacc**. If you intend to use the UNIX utility, your time will be better spent reading the UNIX documentation than the current appendix.

LLama and **occ**s are both discussed in one appendix because they are very similar at the user level. I will discuss the common parts of both systems first, and then describe the individual characteristics of **occ**s and **LLama** in separate sections. Both programs are *compiler compilers*, because they are themselves compilers which translate a very-high-level-language description of a compiler to a high-level-language source code for a compiler. I will use the term *compiler compiler* when I'm talking about characteristics shared by both systems.

E.1 Using The Compiler Compiler

Like **LEX**, **occ**s and **LLama** are C preprocessors. They take as input a file that contains an augmented, attributed grammar that describes a programming language, and they create the C source code for a parser for that language. C source-code fragments that specify code-generation actions in the input are passed to the output with various translations made in that source code to give you access to the parser's value stack. **Occs** outputs the source for a bottom-up parser (both the tables and the parser itself);

LLama outputs a top-down parser. The generated parser subroutine is called `yyparse()`—just call this subroutine to get the parsing started. The subroutine returns zero if it parses the input successfully, `-1` if an unrecoverable error in the input is encountered. A mechanism is provided for you to return other values as well. `yyparse()` expects to get input tokens from a scanner subroutine generated by L^EX, the interaction was discussed at the beginning of Appendix D.

The parser subroutine,
`yyparse()`.

Occs, despite superficial similarities, is not `yacc`—I've made no attempt to perpetuate existing deficiencies in the UNIX program in order to get compatibility with UNIX. **Ococs** has a better error-recovery mechanism than `yacc` (though it's still not ideal), the output code is more maintainable, and it provides you with a considerably improved debugging environment. This last change, activated when you compile with `YYDEBUG` defined, gives you a window-oriented debugging system that lets you actually watch the parse process in action. You can watch the parse stack change, see attributes as they're inherited, set breakpoints on various stack conditions and on reading specified symbols, and so forth. I've documented differences between `occs` and `yacc` in footnotes; these are of interest only if you're already familiar with `yacc`. **LLama**, of course, is radically different from `yacc` in that it uses a completely different parsing technique. Don't be misled by the superficial similarities in the input formats.

Unfortunately, it's difficult to present either compiler compiler in a hierarchical fashion—you may have to read this Appendix twice. Similarly, the description of the attribute-passing mechanism will make more sense after you've seen it used to build the C compiler in Chapter Six.

E.2 The Input File

The **LLama** and `occs` input files, like L^EX input files, are split up into three sections, separated from one another with `%%` directives:

Input-file organization.

definitions

`%%`

rules

`%%`

code

Both the definitions and code sections can be empty.

E.3 The Definitions Section

The first part of the input file contains both directives to the compiler compiler itself and C source code that is used later on in the rules section. As with the L^EX input format, the code must be surrounded by `%{` and `%}` directives.¹ All the directives start with a percent sign, which must be in the leftmost column. Those directives that are unique to one or the other of the programs are discussed later, but those directives that are shared by both **LLama** and `occs` are summarized in Table E.1.

Code in the definitions section.

The `%token` and `%term` directives are used to declare nonterminal symbols in the grammar. The directive is followed by the name of the symbol, which can then be used

Defining tokens:
`%token`, `%term`.

1. Yacc passes lines that start with white space through to the output without modification, as if they were part of a `%{ ... %}` delimited code block. `Ococs` does not support this mechanism.

Table E.1. Occs and LLama % Directives

Directive	Description
%%	Separates the three sections of the input file.
%{	Starts a code block in the definitions section. All lines that follow, up to a line that starts with a %} directive, are written to the output file unmodified.
%}	Ends a code block.
%token	Defines a token.
%term	A synonym for %token.
/* */	C-like comments are recognized—and ignored—by occ , when found outside of a %{...%}-delimited code block. They are passed to the output if found inside a code block.

later in a production in the rules section. For example:

```
%term LP RP /* ( and ) */  
%term ID /* an identifier */  
%term NUM /* a number */  
%token PLUS STAR /* + * */
```

yyout.h: Token definitions

Several names can be listed on a single line. **Occs** creates an output file called *yyout.h* which holds symbolic values for all these tokens, and these values must be returned from the lexical analyzer when the equivalent string is recognized.² For example, the previous declarations produce the following *yyout.h* file:

```
#define _EOI_ 0  
#define LP 1  
#define RP 2  
#define ID 3  
#define NUM 4  
#define PLUS 5  
#define STAR 6
```

Token names.

The arguments to %term are used verbatim for the macro names in *yyout.h*. Consequently, they must obey the normal C naming conventions: The token names must be made up of letters, numbers, and underscores, and the first symbol in the name cannot be a number.³ The definitions in *yyout.h* can then be used from within a LEX input file to build a lexical analyzer. A sample LEX file is in Listing E.1.

Code Blocks %{...%}

Code blocks, delimited by %{ and %} directives, are also found in the definitions section. These blocks should contain the normal things that would be found at the top of a C file: macro definitions, **typedefs**, global-variable declarations, function prototypes, and so forth. Certain macro names are reserved for use by the compiler compiler, however, and these are summarized in Table E.2.⁴ You can modify the parser's behavior by defining these macros yourself in a code block in the definitions section (there's no need to **#undef** them first).

-
2. Yacc creates a file called *y.tab.h*, and the file is created only if *-d* is specified on the command line.
 3. Yacc doesn't impose these restrictions. Also, **yacc** (but not **occ**) accepts a definition of the form %token NAME *number*, in which case it uses the indicated number for the token value rather than assigning an arbitrary value.
 4. Of these, YYABORT, YYACCEPT, YYDEBUG YYMAXDEPTH, and YYSTYPE are supported by **yacc**. The YYDEBUG flag just activates a running log of shift and reduce actions to the screen—there is no interactive debugging environment.

Listing E.1. *expr.lex*—L^EX Input Specification for Simple Expressions

```

1  %{
2      #include "yyout.h"
3  %}
4  %%
5  "+"     return PLUS;
6  "*"     return STAR;
7  "("     return LP;
8  ")"     return RP;
9  [0-9]+  return NUM;
10 [a-z]+ return ID;
11 %%
```

In addition to the foregoing, the file *<stdio.h>* is automatically **#included** at the top of the parser file before any code that you specify in the definitions section is output. It's harmless, though unnecessary, for you to **#include** it explicitly. Similarly, "yyout.h" and the ANSI variable-argument-list-definitions file, *<stdarg.h>*, is also **#included** automatically.⁵ Finally, the yacc version of the stack macros are included (*<tools/yystack.h>*). These are described in Appendix A.

Note that all the symbols that **occs** uses start with YY, yy, or Yy. You should not start any of your own names with these characters to avoid possible name conflicts.

Internal names: yy, YY,
YY

E.4 The Rules Section

The rules section of the input file comprises all lines between an initial %% directive and a second %% directive or end of file. Rules are made up of augmented, attributed productions: grammatical rules and code to be executed. A rules section for an **occs** input file is shown in Listing E.2 along with the necessary definitions section. (I'll discuss this file in greater depth in a moment; it's presented here only so that you can see the structure of an input file.)

Productions take the form:

Occs' modified BNF:
representing productions.

left-hand_side : *right-hand_side* | *rhs_2* | ... | *rhs_N* ;

A colon is used to separate the left- from the right-hand side instead of the → that we've used hitherto. A vertical bar (|) is used to separate right-hand-sides that share a single left-hand side. A semicolon terminates the whole collection of right-hand sides. Terminal symbols in the grammar must have been declared with a previous %term or %token, and their names must obey the C naming conventions.

Nonterminal names can be made up of any collection of printing nonwhite characters except for the following:

Nonterminal names.

% { } [] () < > * : | ; ,

There can't be any imbedded spaces or tabs. Imbedded underscores are permitted. Imbedded dots can be used as well, but I don't suggest doing so because dots make the *yyout.doc* file somewhat confusing.⁶ I'd suggest using lower-case letters and underscores

5. Yacc actually puts the token definitions in the output file. Occs puts the definitions in *yyout.h* and **#includes** it in the output file.

6. Exactly the opposite situation applies to yacc. An imbedded underscore should be avoided because an underscore is used to mark the current input position in *yyout.doc*.

Table E.2. Macros that Modify the Parser

YYABORT	This macro determines the action taken by the parser when an unrecoverable error is encountered or when too many errors are found in the input. The default macro is: <pre>#define YYABORT return(1)</pre> but you can redefine it to something else in a code block at the top of the <code>occ</code> s input file if you prefer. Note that there are no parentheses attached to the macro-name component of the definition.
YYACCEPT	This macro is like YYABORT, except that it determines the accepting action of the parser. The standard action is <code>return 0</code> , but you can redefine it to something else if you prefer.
YYCASCADE	(Occs only.) This constant controls the number of parser operations that have to follow an error before another error message is printed. It prevents cascading error messages. It's initialized to five. #define it to a larger number if too many error messages are printed, a smaller number if too many legitimate error messages suppressed.
YYD(x)	If the <code>-D</code> switch is specified, then all text in a YYD macro is expanded, otherwise it's ignored. For example, the following <code>printf()</code> statement is executed only if <code>-D</code> is present on the <code>occ</code> s command line. <pre>YYD(printf("hi there");)</pre>
YYDEBUG	Set automatically by the <code>-D</code> command-line switch to enable the interactive debugging environment, but you can define it explicitly if you like. The contents are unimportant.
YYMAXDEPTH	This macro determines the size of the state and value stacks. It's 128 by default, but you may have to enlarge it if you get a "Parse stack overflow" error message at run time. You can also make it smaller for some grammars.
YYMAXERR	(Occs only.) Initially defined to 25, the value of this macro determines when the parser will quit because of too many syntax errors in the input.
YYPRIVATE	Defining this macro as an empty macro makes various global-variable declarations public so that you can see them in a link map. Use the following: <pre>#define YYPRIVATE /*nothing*/</pre>
YYSHIFTACT(tos)	(Occs only.) This macro determines the default shift action taken by the parser. That is, it controls what gets put onto the value stack when an input symbol is shifted onto the stack. It is passed a pointer to the top of the value stack after the push (<code>tos</code>). The default definition looks like this: <pre>#define YYSHIFTACT(tos) ((tos)[0]=yyval)</pre> (<code>yyval</code> is used by the lexical analyzer to shift an attribute associated with the current token. It is discussed below.) You can redefine YYSHIFTACT as an empty macro if none of the tokens have attributes.
YYSTYPE	This macro determines the typing of the value stack. It's discussed at length in the text.
YYVERBOSE	If this macro exists, the parser uses more verbose diagnostic messages when YYDEBUG is also defined.

Listing E.2. *expr.y*—Occs Rules Section for the Expression Compiler

```

1 %term ID          /* an identifier      */
2 %term NUM         /* a number           */
3 %left PLUS        /* +                 */
4 %left STAR        /* *                 */
5 %left LP RP       /* ( and )           */
6 %%
7 /* A small expression grammar that recognizes numbers, names, addition (+),
8 * multiplication (*), and parentheses. Expressions associate left to right
9 * unless parentheses force it to go otherwise. * is higher precedence than +.
10 * Note that an underscore is appended to identifiers so that they won't be
11 * confused with rvalues.
12 */
13
14 s      : e ;
15
16 e      : e PLUS e { yycode("%s += %s\n", $1, $3); free_name( $3 ); }
17 | e STAR e { yycode("%s *= %s\n", $1, $3); free_name( $3 ); }
18 | LP e RP { $$ = $2;
19 | NUM       { yycode("%s = %s\n", $$ = new_name(), yytext ); }
20 | ID        { yycode("%s = _%s\n", $$ = new_name(), yytext ); }
21 ;
22 %%

```

for nonterminal names and upper-case letters for tokens, as in Listing E.2. This way you can tell what something is by looking at its name. It's a hard error to use a nonterminal that is not found somewhere to the left of a colon.

Text surrounded by curly braces is code that is executed as a parse progress. In a top-down **LLama** parser, the code is executed when it appears at the top of the parse stack. In the bottom-up **occs** parser, the code is executed when a reduction by the associated production is performed.

ϵ productions are represented with an empty right-hand side (and an optional action). It's best to comment the fact that the production has been left empty on purpose. All the following are legitimate ϵ productions:

```

a : /* empty */
;

b : /* epsilon */ { an_action(); }

c : normal_rhs   { an_action(); }
| /* epsilon */ { another_action(); }
;

```

I'd suggest using the formatting style shown in Listing E.2. The colons, bars, and semicolons should all line up. The nonterminal name should start at the left column and the body of a production should be indented. The curly braces for the actions all line up in a nice column. If a name is too long, use the following:

```

very_long_left_hand_side_name
: right_hand_side    { action(1); }
| another            { action(2); }
;

```

Any valid C statement that's permitted inside a C subroutine is allowed in an action, with the following caveats:

Augmentations: putting code in a production.

ϵ productions.

Formatting.

return and **break** statements in actions.

- (1) A **break** statement at the outermost level terminates processing of the current action but not the parser itself.
- (2) All **return** statements must have an associated value. A simple **return** is not permitted.
- (3) A **return** statement causes the parser to return only if it takes a nonzero argument. That is, **return** 0 causes processing of the current action to stop, but the parsing continues; **return** 1 forces the parser to return to the calling routine, because the argument is nonzero. The argument to **return** is returned from `yyparse()`.
- (4) The YYACCEPT and YYABORT macros determine what happens when the parser terminates successfully (accepts the input sentence) or fails (rejects the input sentence). The default actions return 0 and 1 respectively, so you should not return either of these values from within an action.

Local variables in actions.

You can't declare subroutines in an action because the actions themselves are inside a subroutine. You can, however, declare local variables after any open brace, and you can use an action that's several lines long. The following is okay:

```
left_hand_side : right_hand_side
{
    int i, monster;
    for( i = 10; --i >= 0 ; )
        bug( i, monster );
}
```

Of course, you can't declare a local variable in one action and access that variable from another. True global variables that are declared in a `%{ ... %}` block in the definitions sections can be accessed everywhere.

The start production.

The first production in the rules section is taken as the start production and is treated specially: Its left-hand side is the goal symbol and it can have only one right-hand side.⁷

E.5 The Code Section

The third part of the input file, the *code section*, comprises all text after the second `%%` directive. This entire section is written to the output file verbatim. It is used the same way as the LEX code section—it should contain subroutines that are called from the actions. Note that the **occ**s code section differs from both the LEX and LLama sections in that dollar attributes (discussed below) can be used in the code found there.⁸

E.6 Output Files

LLama and occs create several output files—their interaction with LEX is discussed at the beginning of Appendix D. The two programs use different names for these files, however, and these are summarized in Table E.3.

7. Some versions of yacc support a `%start` directive that lets you define an arbitrary production as the start production. This mechanism isn't supported by occs.

8. Yacc doesn't permit this: dollar attributes may be used only in the rules sections.

Table E.3. Files Generated by **Occs** and **LLama**

Occs Name	LLama Name	Command-line Switches	Contains
yyout.c	llout.c	<i>none</i> , -p	Normally contains parser and action code taken from original grammar. If <i>-p</i> is specified to occ s, this file contains the parser only. (LLama doesn't support a <i>-p</i> switch.)
yyact.c		-a	If <i>-a</i> is specified to occ s, this file contains code for the action components of the productions and all of the code portion of input file (but not the parser). LLama doesn't support <i>-a</i> .
yyout.h yyout.sym	llout.h llout.sym	<i>none</i> -s, -S, -D	Contains #defines for the tokens. The symbol table. The <i>-S</i> version contains more information than the <i>-s</i> version. Specifying <i>-D</i> implies <i>-s</i> —you can get more-verbose tables with <i>-DS</i> .
yyout.doc		-v, -V	Used only by occ s, contains the LALR(1) state-machine description along with any warning messages that are sent to the screen.

E.7 Command-Line Switches

The basic command-line usage for both programs is the same:

```
occ [−switches] file
llama [−switches] file
```

where the *file* is an input specification file, as just discussed. The programs take the following command-line switches:⁹

- −a (**occ**s only.) Rather than creating a combined parser-and-action file, the action subroutine only is output to the file *yyact.c*. The parser component of the output file can be created using the *-p* switch, described below. Generate Action Subroutine.
- −c[N] (**LLama** only.) This switch controls the parse-table compression method; it works just like the **LEX** switch with the same name. It changes the compression algorithm to one in which each character/next-state pair is stored as a two-byte object. This method typically (though not always) gives you smaller tables and a slower parser than the default method. The *-c* switch takes an optional numeric argument (*-c5*, for example) that specifies the threshold above which the pair compression kicks in for a given row of the table. If the row has more than the indicated number of nonerror transitions, it is not compressed. The default threshold, if no number is specified, is four. Pair compress **LLama** parse tables.
- −D Enables the interactive debugging environment in the generated parser. All that this switch really does is insert a **#define YYDEBUG** into the output file. You can do the same thing yourself by putting an explicit **#define** in the input specification. Specifying *-D* also causes *-s* to be set. Activating the interactive debugging environment.
- −f (**LLama** only.) This switch causes the output tables to be uncompressed, thereby speeding up the parser a little at the cost of program size. Uncompressed tables.

9. Of these, only *-v* is supported by **yacc**, though **yacc** is considerably less verbose than **occ**s. **occ**s' *-D* switch is nothing like **yacc**'s *-d* switch.

Making private variables public.	-g	Several global variables and subroutines are used in the output parser. These are declared <code>static</code> by default, so their scope is limited to <code>yyout.c</code> (or <code>llout.c</code>). The -g switch causes the compiler compiler to make these variables global to the entire program. You can achieve the same effect by putting the following in the first section of the input specification:
		<code>%{ #define YYPRIVATE %}</code>
Suppress <code>#line</code> directives.	-l	(<i>I</i> is an ell.) Occs and LLama usually generate <code>#line</code> directives that cause the C compiler's error messages to reflect line numbers in the original input file rather than the output file. (A preprocessor directive of the form <code>#line N "file"</code> tells the compiler to pretend that it's on line <i>N</i> of the indicated <i>file</i> .) These directives can cause problems with source-level debuggers, however, and the -l switch gives you a way to turn them off. Generally you should not use -l until you've eliminated all the syntax errors from the input file. Resist the temptation to modify the output file in order to fix a syntax error. It's too easy to forget to modify the input too, and the output file is destroyed the next time you use the compiler compiler.
Alternative parser templates.	-m	Occs and LLama both generate only part of the parser. They create the parse tables and a few macros and <code>typedefs</code> needed to compile those tables. The remainder of the parser is read from a template file, called <i>occsp.par</i> (in the case of occsp) and <i>llamap.par</i> (in the case of LLama). The programs look for these files, first in the current directory, and then along all directories specified in the LIB environment variable. LIB should contain a semicolon-delimited list of directories. The -m command-line switch is used to specify an alternate template file. For example:
		<code>occsp -m/usr/allen/template/template.par</code>
The LIB environment variable.		tells occsp to read a file called <i>template.par</i> from the <i>/usr/allen/template</i> directory rather than using the default template. The template files are discussed in depth in Chapters Four and Five. Note that there's no space between the <i>m</i> and the first character of the path name.
Create parser only.	-p	(occsp only.) Rather than creating a combined parser-and-action file, the parser component of the file is output to the file <code>yyout.c</code> . No actions are put into <code>yyout.c</code> —the action component of the output file can be created using the -a switch, described earlier.
Symbol Tables.	-S	These switches causes the symbol-table file (<code>yyout.sym</code> for occsp and <code>llout.sym</code> for LLama) to be generated. A capital S causes the symbol-table file to contain more information—the FIRST() sets for each nonterminal are printed in the occsp table, and both FIRST and FOLLOW sets are printed in LLama 's.
Use stdout.	-t	Causes occsp to output the parser to standard output rather than to the <code>yyout.c</code> file.
Send large tables to <code>yyoutab.c</code> .	-T	The table-compression method used by occsp can create several hundred small arrays, and some compilers can't handle that many declarations in a single file. The -T switch splits the parser file up into two portions, one containing the parser and most of the tables, the other containing the two biggest tables—the ones most likely to give the compiler problems. These parser is put into <code>yyout.c</code> as usual. The two parse tables are put into <code>yyoutab.c</code> . The two files can then be compiled independently and linked together. You can use this switch in conjunction to the -a and -p switches to split the parser into three parts. Use:

<code>occ</code> s -pT input.y	<i>Create yyout.c and yyparse.c</i>
<code>occ</code> s -a input.y	<i>Create actions file</i>

- V The `-v` switch puts `occ`s and **LLama** into verbose mode. `occ`s generates a file called `yyout.doc` which contains a human-readable description of the LALR(1) state machine. There is no equivalent file from **LLama**. Both programs send progress reports to the screen as they work and they print various statistics describing table sizes, and so forth, when done. `-V` works like `-v` except that more information about the internal workings of `occ`s and **LLama** than you're interested in seeing is printed on the screen. It's used primarily for debugging `occ`s itself or for seeing how the LALR(1) state machine is put together. `-V` also puts the LALR(1) lookaheads for every kernel item (rather than just for the items that cause reductions) into `yyout.doc`.
- w Suppress all warning messages. These warnings announce the presence of shift/reduce and reduce/reduce conflicts found when the state machine is constructed. Once you've ascertained that these warnings are harmless, you can use `-w` to clean up the screen. Warnings are printed to `yyout.doc`, even if `-w` is specified, and the number of warnings is always printed at the end of an `occ`s run, even if the warnings themselves aren't printed.
- W Normally, the number of hard errors is returned to the operating system as the exit status. If `-W` is specified, the exit status is the sum of the number of hard errors and warning messages. This switch is handy if you want the UNIX **make** utility (or equivalent) to terminate on warnings as well as errors.

Verbose-mode LALR(1)
state-machine descrip-
tion.

SUPPRESS warning mes-
sages.

E.8 The Visible Parser¹⁰

An `occ`s- and **LLama**-generated parser can function in one of two ways. In *production mode*, it just parses the input and generates code, as you would expect. In *debug mode*, however, a multiple-window interactive debugging environment (or IDE) is made available to you. Debug mode is enabled in one of two ways, either specify `-D` on the `occ`s command-line, or put the following in the rules section of your input file:

```
%{ #define YYDEBUG %}
```

The resulting file should be compiled and linked to three libraries: *l.lib*, *curses.lib*, and *termlib.lib*.¹¹ Only the first of these libraries is necessary if you're not compiling for debugging. The `occ`s source file for a small expression compiler is on the software-distribution disk mentioned in the preface, along with the curses and termlib libraries. You may need to recompile the libraries to get them to work properly if you're not using Microsoft C (or QuickC). An executable version of the compiler (*expr.exe*) is also on the disk, however, so you don't have to recompile right now. Make the expression compiler with Microsoft C as follows:

Production and debug
mode.

IDE (interactive debug-
ging environment).

Compiling for Debugging.

10. This section describes how to use the interactive debugging environment. If at all possible, you should have an `occ`s output file, compiled for interactive debugging, running in front of you as you read. An executable version of the expression compiler described below is provided on the software distribution disk in the file *expr.exe*.

11. Use cc files... -l curses -l termlib in UNIX.

```
occg -vD expr.y
lex -v expr.lex
cl -o expr.exe yyout.c lexyy.c -link l.lib curses.lib termlib.lib
```

Then run it with:

```
expr test
```

where *test* is an input file containing a simple expression composed of numbers, variable names, plus signs, asterisks, and parentheses. The expression must be the first and only line in the file.

Direct-video output.

The VIDEO environment variable.

Window names.

Changing stack-window size from the command line.

In the MS-DOS environment, the debugger uses the ROM-BIOS for its output. It can also use direct-video writes, however, and on some clones this approach is more reliable than using the BIOS. To activate the direct-video routines, set the VIDEO environment to DIRECT with the following command to COMMAND.COM:

```
set VIDEO=DIRECT
```

The direct-video routines assume that either an MGA, or a CGA running in high resolution, 25×80 text mode is present (it automatically selects between these two display adapters). Most other video adapters can simulate one or the other of these modes. The program will not work with a CGA running in one of the other text modes (25×20 or 25×40), however. Issue a MODE BW80 request at the DOS prompt to get the CGA into high-resolution, text mode.

The program comes up with several empty windows, as shown in Figure E.1. The *stack* window is used to show the state and value-stack contents; the *comments* window holds a running description of the parse process; generated code is displayed in the *output* window, the *lookahead* window displays the most-recently read token and lexeme—the current lookahead symbol, and the *prompts* window is used to communicate with the debugger. The size of the stack window can be changed by specifying an optional *-sN* command-line switch when you invoke the program, as in *expr -s15 test*. In this case the stack window occupies 15 lines. The other windows scale themselves to take up what's left of the screen.

Several commands are supported by the debugger; you can see a list by typing a question mark, which prints the list shown in Figure E.2. The easiest command to use is the *g* command (for go) which just starts up the parsing. Figure E.3 shows the various windows several stages into parsing the expression *a*(b+c)*. The parsing can be stopped at any time by pressing any key.

In Figure E.3, the parser has just read the *c* (but it has not yet been shifted onto the parse stack). The ID *c* in the *lookahead* window is the most recently read token and lexeme. The *output* window shows code that has been output so far, and the current parse stack is shown in the *stack* window.

The *output* window shows output sent to all three segment streams. The leftmost column identifies the segment—it holds a C for code, D for data, or B for bss. Because the window is so narrow, sequences of white-space characters (tabs, spaces, and so forth) are all replaced with a single space character in this window.

The *stack* window always displays the topmost stack elements. If the stack grows larger than the window, the stack window scrolls so that the top few stack elements are displayed. The three columns in the *stack* window are (from left to right): the parse stack, the parse stack represented in symbolic form rather than numbers, and the value stack. Several commands other than *g* are recognized at the "Enter command" prompt in the *prompts* window:

Figure E.1. The Initial Debug Screen

```
+-----[stack]-----+
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
+-----[comments]-----+[output]-----+
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
| | | | | |
+-----[prompts]-----+[lookahead]---+
| Enter command (space to continue, ? for list): | |
+-----+-----+
```

Figure E.2. Debug-Mode Commands

a (A)bort parser by reading EOI	l (L)og output to file
b modify or examine (B)reakpoint	n/N (N)oninteractive mode
d set (D)elay time for go mode	q (Q)uit (exit to dos)
f read (F)ile	r (R)efresh stack window
g (G)o (any key stops parse)	w (W)rite screen to file or device
i change (I)nput file	x Show current and prev. le(X)eme

Space or Enter to single step

space, enter You can single step through the parse by repetitively pressing the space bar or the Enter key. In practice, this is usually more useful than *g* which can go by faster than you can see. Single step.

Ctrl-A, Ctrl-B These two commands are deliberately not printed on the help screen. They are debugger hooks. Ctrl-A causes the debugger to call the subroutine *yyhook_a()*, and Ctrl-B calls *yyhook_b()*. The default versions of these routines do nothing. You can use these commands for two things. First, its often useful to run the IDE under a compiler debugger like *sdb* or *CodeView*, and you can use a hook as an entry point to the debugger: Set a breakpoint at *yyhook_a()* and then start the parser running; issuing a Ctrl-A then returns you to the debugger. Debugger hooks.

Figure E.3. The Windows Several Steps into a Parse

```
+-----[stack]-----+
| | | | |
| 6| PLUS | _t1
| 5|   | _t1
| 3| LP | _t0
| 7| STAR | _t0
| 4|   | _t0
| 0| S | <empty>
+-----[comments]-----+[output]-----+
|Shift <LP> (goto 3)          |B| public word _t0,_t1,_t2,_t3;
|Advance past ID <b>          |B| public word _t4,_t5,_t6,_t7;
|Shift <ID> (goto 1)          |C| _t0 = a
|Advance past PLUS <+>       |C| _t1 = b
|Reduce by (5) e->ID          |
|  (goto 5)                   |
|Shift <PLUS> (goto 6)         |
|Advance past <PLUS>          |
|Read ID <c>                 |
+-----[prompts]-----+[lookahead]---+
| Enter command (space to continue, ? for list): | ID c |
+-----+
```

The two hook subroutines are in modules by themselves, so you can use them to add capabilities to the IDE by linking your own versions rather than the library versions. For example, it's occasionally useful to print the symbol table when you are debugging the parts of the compiler that are processing declarations. This capability can't be built into the IDE because there's no way of knowing what the data structures will look like. The problem is solved by supplying a subroutine called `yyhook_a()` which prints the symbol table. Your version is linked rather than the default library version, so a Ctrl-A command now prints the symbol table.

- a** This command aborts the parse process by forcing the lexical analyzer to return the end-of-input marker on all subsequent calls. Compare it with the `q` command, which exits immediately to the operating system without returning to the parser itself.
- Breakpoints.** **b** The `b` command is used to set, clear, or examine breakpoints. Four types of breakpoints are supported:
 - `input` (i) breaks when a specified symbol is input.
 - `input line` (l) breaks when the input sweeps past the specified input line. This breakpoint is activated only when the lexical analyzer returns a token, and it's possible to specify an input line that doesn't have any tokens on it. (In this case, the break occurs when the first token on a line following the indicated one is read.)

Note that this breakpoint, unlike any of the others, is automatically cleared as soon as it is triggered.

stack (s) breaks as soon as a specified symbol appears at the top of the parse stack. In a **LLama** parser, parsing stops before any action code is executed.

production (p) breaks just before the parser applies a specific production. In **occs**, it breaks just before the reduction. In **LLama**, it breaks just before replacing a nonterminal with its right-hand side.

When you enter *b*, you are prompted for one of these types (enter an *i*, *l*, *s*, or *p* for “input,” “line”, “stack,” or “production”). You are then prompted for a symbol. The different kinds of breakpoints can take different inputs in response to this prompt:

input (1) A decimal number is assumed to be a token value, as defined in *yyout.h* or *yyout.sym* (*llout.h* or *llout.sym* if **LLama**). Parsing stops when that token is read. (2) A string that matches a token name (as it appears in a `%term` or `%token` directive) behaves like a number—parsing stops when that token is read. (3) Any other string causes a break when the current lexeme matches the indicated string. For example, if you have a `%term NAME` in the input file and the compiler compiler puts a `#define NAME 1` in *yyout.h* or *llout.h*, then entering either the string 1 or the string NAME causes a break the next time a NAME token is read. You can also enter the string Illiavitch, whereupon a break occurs the next time *Illiavitch* shows up as a lexeme, regardless of the token value.

input line Enter the input line number.

stack (1) If you enter a number, parsing stops the next time that number appears at the top of the real parse stack (the leftmost column in the *stack* window). (2) Any other string is treated as a symbolic name, and parsing stops the next time that name appears at the top of the symbol stack (the middle column in the window).

production You must enter a number for this kind of breakpoint. Parsing stops just before the particular production is applied (when it's replaced in a **LLama** parser and when a reduction by that production occurs in an **occs** parser). Production numbers can be found in the *yyout.sym* or *llout.sym* file, generated with a `-s` or `-S` command-line switch.

There's no error checking to see if the string entered for a stack breakpoint actually matches a real symbol. The breakpoint-processing routines just check the input string against the symbol displayed in the symbolic portion of the stack window. If the two strings match, the parsing stops. The same technique is used for input breakpoints. If the symbol isn't a digit, then the breakpoint string is compared, first against the input lexeme, and then by looking up the name of the current lexeme. If either string matches, then parsing stops.

Three other breakpoint commands are provided: *c* clears all existing breakpoints, *d* displays all existing breakpoints, and *?* prints the help screen shown in Figure E.4.

d The *d* command is used to slow down the parse. The parsing process, when started with a *g* command, nips along at a pretty good clip, often too fast for you to see what's going on. The *d* command is used to insert a delay between

Adding a delay in go mode.

Figure E.4. Breakpoint Help Screen

```
Select a breakpoint type (i,p,or s) or command (c or l)

Type: Description:           Enter breakpoint as follows:
      i     input.....number for token value
            or string for lexeme or token name
      l     input line.....input line number
      p     reduce by production.....number for production number
      s     top-of-stack symbol.....number for state-stack item
            or string for symbol-stack item
      s = clear all breakpoints
      l = list all breakpoints
```

every parse step. Setting the delay to 0 puts it back to its original blinding speed. Delay times are entered in seconds, and you can use decimal fractions (1, 2.5, .5, and so forth) if you like.

- | | | |
|---|-----------------------|--|
| Examine file. | f | The <i>f</i> command lets you examine a file without leaving the debugger. You are prompted for a file name, and the file is then displayed in the stack window, one screenfull at a time. |
| Go! | g | The <i>g</i> command starts the parse going. |
| Specify input file. | i | The <i>i</i> command lets you change the input file from the one specified on the command line. It prompts you for a file name. |
| Log all output to file. | l | This command causes the entire parse to be logged to a specified file, so that you can look at it later. If you're running under MS-DOS, you can log to the printer by specifying <i>prn:</i> instead of a file name. Some sample output is shown in Figures E.5 and E.6. Output to the CODE window is all preceded by CODE->. Most other text was COMMENT-window output. The parse stack is drawn after every modification in one of two ways (you'll be prompted for a method when you open the log file). Listing E.3 shows vertical stacks and Listing E.4 shows horizontal stacks. The latter is useful if you have relatively small stacks or relatively wide paper—it generates more-compact log files in these situations. The horizontal stacks are printed so that items at equivalent positions on the different stacks (parse/state, symbol, and value) are printed one atop the other, so the column width is controlled by the stack that requires the widest string to print its contents—usually the value stack. If you specify horizontal stacks at the prompt, you will be asked which of the three stacks to print. You can use this mechanism to leave out one or more of the three stacks. |
| Specifying horizontal stacks. | n,
N | These commands put the parser into noninteractive mode. The n command generates a log file quickly, without having to watch the whole parse happen before your eyes. All screen updating is suppressed and the parse goes on at much higher speed than normal. A log file must be active when you use this command. If one isn't, you'll be prompted for a file name. You can't get back into normal mode once this process is started. The N command runs in noninteractive mode without logging anything. It's handy if you just want to run the parser to get an output file and aren't interested in looking at the parse process. |
| Noninteractive mode:
Create log without window updates | n,
N | |
| Run parser without logging or window updates. | | |

Listing E.3. Logged Output—Vertical Stacks

```

1 CODE->public word _t0,_t1,_t2,_t3;
2 CODE->public word _t4,_t5,_t6,_t7;
3
4 Shift start state
5 +-----+
6 0| 0|           S | -
7 +-----+
8 Advance past NUM <1>
9 Shift <NUM> (goto 2)
10 +-----+
11 0| 2|           NUM | -
12 1| 0|           S | -
13 +-----+
14 Advance past PLUS <+>
15 CODE->_t0 = 1
16
17 Reduce by (4) e->NUM
18 +-----+
19 0| 0|           S | -
20 +-----+
21 (goto 4)
22 +-----+
23 0| 4|           e | _t0
24 1| 0|           S | -
25 +-----+
26 Shift <PLUS> (goto 6)
27 +-----+
28 0| 6|           PLUS | _t0
29 1| 4|           e | _t0
30 2| 0|           S | -
31 +-----+
32 Advance past NUM <2>
33 Shift <NUM> (goto 2)
34 +-----+
35 0| 2|           NUM | _t0
36 1| 6|           PLUS | _t0
37 2| 4|           e | _t0
38 3| 0|           S | -
39 +-----+
40 Advance past STAR <*>
41 CODE->_t1 = 2

```

- | | |
|--|--|
| <p>q returns you to the operating system.</p> <p>r An <i>r</i> forces a STACK-window refresh. The screen-update logic normally changes only those parts of the STACK window that it thinks should be modified. For example, when you do a push, the parser writes a new line into the stack window, but it doesn't redraw the already existing lines that represent items already on the stack. Occasionally, your value stack can become corrupted by a bug in your own code, however, and the default update strategy won't show you this problem because it might not update the incorrectly modified value stack item. The <i>r</i> command forces a redraw so that you can see what's really on the stack.</p> <p>w (for write) dumps the screen to an indicated file or device, in a manner analogous to the Shift-PrtSc key on an IBM PC. This way you can save a snapshot</p> | <p>Quit.</p> <p>Redraw stack window.</p> <p>Save screen to file.</p> |
|--|--|

Listing E.4. Logged Output—Horizontal Stacks

```

1 CODE->public word _t0,_t1,_t2,_t3;
2 CODE->public word _t4,_t5,_t6,_t7;
3
4 Shift start state
5     PARSE 0
6     SYMBOL S
7     ATTRIB -
8 Advance past NUM <1>
9 Shift <NUM> (goto 2)
10    PARSE 0 2
11    SYMBOL S NUM
12    ATTRIB - -
13 Advance past PLUS <+>
14 CODE->_t0 = 1
15
16 Reduce by (4) e->NUM
17     PARSE 0
18     SYMBOL S
19     ATTRIB -
20     (goto 4)
21     PARSE 0 4
22     SYMBOL S e
23     ATTRIB - _t0 (e)
24 Shift <PLUS> (goto 6)
25     PARSE 0 4 6
26     SYMBOL S e PLUS
27     ATTRIB - _t0 (e) _t0 (PLUS)
28 Advance past NUM <2>
29 Shift <NUM> (goto 2)
30     PARSE 0 4 6 2
31     SYMBOL S e PLUS NUM
32     ATTRIB - _t0 _t0 _t0
33 Advance past STAR <*>
34 CODE->_t1 = 2

```

of the current screen without having to enable logging. Any IBM box-drawing characters used for the window borders are mapped to dashes and vertical bars. I used the *w* command to output the earlier figures. Note that the output screen is truncated to 79 characters because some printers automatically print a linefeed after the 80th character. This means that the right edge of the box will be missing (I put it back in with my editor when I made the figures).

Display lexeme.

x

Display both the current and previous lexeme in the comments window. The token associated with the current lexeme is always displayed in the *tokens* window.

E.9 Useful Subroutines and Variables

There are several useful subroutines and variables available in an **occ**- or **LLama**-generated parser. These are summarized in Table E.4 and are discussed in this section.

void yyparse()

The parser subroutine.

This subroutine is the parser generated by both **occS** and **LLama**. Just call it to get the parse started.

char *yypstk(YYSTYPE *val, **char** *sym)

Print value-stack item.

This subroutine is called from the debugging environment and is used to print the value stack. It's passed two pointers. The first is a pointer to a stack item. So, if your value stack is a stack of character pointers, the first argument will be a pointer to a character pointer. The second argument is always a pointer to a string holding the symbol name. That is, it contains the symbol stack item that corresponds to the value stack item. The returned string is truncated to 50 characters, and it should not contain any newline characters. The default routine in *l.lib* assumes that the value stack is the default **int** type. It's shown in Listing E.5. This subroutine is used in slightly different ways by **occS** and **LLama**, so is discussed further, below.

int yy_get_args(**int** argc, **char** **argv)

Change stack-window size and specify input file to debugger.

This routine can be used to modify the size of the *stack* window and to open an input file for debugging. The other windows automatically scale as appropriate, and the stack window is not allowed to get so large that the other windows disappear. Typically, *yy_get_args()* is called at the top of your *main()* routine, before *yyparse()* is called. The subroutine is passed *argv* and *argc* and it scans through the former looking for an argument of the form *-sN*, where *N* is the desired stack-window size. The first argument that doesn't begin with a minus sign is taken to be the input file name. That name is not removed from *argv*. All other arguments are ignored and are not removed from *argv*, so you can process them in your own program. The routine prints an error message and terminates the program if it can't open the specified input file. Command-line processing stops immediately after the input file name is processed. So, given the line:

```
program -x -s15 -y foo -s1 bar
```

argv is modified to:

```
program -x -y foo -s1 bar
```

the file *foo* is opened for input, and the stack window will be 15 lines high. A new value of *argc* that reflects the removed argument is returned.

This routine can also be used directly, rather than as a command-line processor. For example, the following sets up a 17-line stack window and opens *testfile* as the input file:

```
char *vects[] = {"", "-s17", "testfile"};
yy_get_args( 3, vects );
```

Using *yy_get_args()* directly.**void** yy_init_occS (YYSTYPE *tos)

Initialization functions.

These routines are called by *yyparse()* after it has initialized the various stacks, but before it has read the first input token. That is, the initial start symbol has been pushed onto the state stack, and garbage has been pushed onto the corresponding entry on the value stack, but no tokens have been read. The subroutine is passed a pointer to the (only) item on the value stack. You can use *yy_init_occS()* to provide a valid attribute for this first value-stack element. A user-supplied version of both functions is also useful when *main()* isn't in the **occS** input file, because it can be used to initialize **static** global variables in the

Table E.4. Useful Subroutines and Variables

int	yyparse	(void);
int	yylex	(void);
char	*yypstk	(void *value_stack_item, char *symbol_stack_item);
void	yycomment	(char *fmt, ...);
void	yycode	(char *fmt, ...);
void	yydata	(char *fmt, ...);
void	yybss	(char *fmt, ...);
int	yy_get_args	(int argc, char **argv);
void	yy_init_occs	(void);
void	yy_init_llama	(void);
void	yy_init_lex	(void);
void	yyerror	(char *fmt, ...);
FILE	*yyout	= stdout; /* output stream for code */
FILE	*yybssout	= stdout; /* output stream for bss */
FILE	*yydataout	= stdout; /* output stream for data */

Listing E.5. *yypstk.c*—Print Default Value Stack

```

1  /* Default routine to print user-supplied portion of the value stack. */
2
3  char    *yypstk(val,sym)
4  void    *val;
5  char    *sym;
6  {
7      static char buf[32];
8      sprintf( buf, "%d", *(int *)val );
9      return buf;
10 }
```

parser file. It's easy to forget to call an initialization routine if it has to be called from a second file. The default subroutines, in *l.lib*, do nothing. (They are shown in Listings E.6 and E.7.)

Listing E.6. *yyinitox.c*—Occs User-Initialization Subroutine

```
1  void yy_init_ox( tos ) void *tos; { }
```

Print parser error messages. **void** yyerror(**char** *fmt,...)

This routine is called from *yyparse()* when it encounters an error, and you should use it yourself for error messages. It works like *printf()*, but it sends output to *stderr* and it adds the current input line number and token name to the front of the message, like this:

Listing E.7. *yyinitll.c*—Default LLama Initialization Function

```
1 void yy_init_llama( tos ) void *tos; { }
```

ERROR (line 00 near TOK) : *your message goes here.*

The *00* is replaced by the current line number and *TOK* is replaced by the symbolic name of the current token (as defined in the *occs* input file). The routine adds a newline at the end of the line for you.

```
void yycomment( char *fmt, ... )
void yycode   ( char *fmt, ... )
void yydata   ( char *fmt, ... )
void yybss    ( char *fmt, ... )
```

Output functions.

These four subroutines should be used for all your output. They work like *printf()*, but write to appropriate windows when the debugging environment is enabled. When the IDE is not active, *yycomment* writes to standard output and is used for comments to the user, *yycode()* writes to a stream called *yycodeout* and should be used for all code, *yydata()* writes to a stream called *yydataout* and should be used for all initialized data, *yybss()* writes to a stream called *ybssout* and should be used for uninitialized data.

All of these streams are initialized to *stdout*, but you may change them at any time with an *fopen()* call. Don't use *freopen()* for this purpose, or you'll close *stdout*. If any of these streams are changed to reference a file the debugger sends output both to the file and to the appropriate window. If you forget and use one of the normal output routines like *puts()* or *printf()*, the windows will get messed up, but nothing serious will happen. *printf()* is automatically mapped to *yycode* if debugging is enabled, so you can use *printf()* calls in the *occs* input file without difficulty. Using it elsewhere in your program causes problems, however.

```
void yyprompt( char *prompt, char *buf, int get_str )
```

Print a message to the prompt window.

This subroutine is actually part of the debugger itself, but is occasionally useful when implementing a debugger hook, described earlier. It prints the prompt string in the IDE's *prompts* window, and then reads a string from the keyboard into *buf*. (there's no boundary checking so be careful about the array size). If *get_str* is true, an entire string is read in a manner similar to *gets()*, otherwise only one character is read. If an ESC is encountered, the routine returns 0 immediately, otherwise 1 is returned.

E.10 Using Your Own Lexical Analyzer

Occs and **LLama** are designed to work with a **LEX**-generated lexical analyzer. You can build a lexical analyzer by hand, however, provided that it duplicates **LEX**'s interface to the parser.

You must provide the following subroutines and variables to use the interactive debugging environment without a lex-generated lexical analyzer:

char *yytext;	<i>current lexeme</i>
int yylineno;	<i>current input line number</i>
int yyleng;	<i>number of characters in yytext[]</i>
int yylex(void);	<i>return next token and advance input</i>
char *ii_ptext(void);	<i>return pointer to previous lexeme</i>
int ii_plength(void);	<i>return length of previous lexeme</i>
int ii_mark_prev(void);	<i>copy current lexeme to previous lexeme</i>
int ii_newfile(char *name);	<i>open new input file</i>

The scanner must be called `yylex()`. It must return either a token defined in `yyout.h` or zero at end of file. You'll also need to provide a pointer to the current lexeme called `yytext` and the current input line number in an `int` called `yylineno`. The `ii_ptext()` subroutine must return a pointer to the lexeme that was read immediately before the current one—the current one is in `yytext` so no special routine is needed to get at it. The string returned from `ii_ptext()` does not have to be '\0' terminated. Like `yyleng`, `ii_plength()` should evaluate the the number of valid characters in the string returned from `ii_ptext()`. `ii_mark_prev()` should copy the current lexeme into the previous one. `ii_newfile()` is called when the program starts to open the input file. It is passed the file name. The real `ii_newfile()` returns a file handle that is used by `yylex()` in turn. Your version of `ii_newfile()` need only do whatever is necessary to open a new input file. It should return a number other than -1 on success, -1 on error. Note that input must come from a file. The debugging routines will get very confused if you try to use `stdin`.

E.11 Occs

This section describes `occ`s-specific parts of the compiler compiler. A discussion of the `LLama`-specific functions starts in Section E.12

E.11.1 Using Ambiguous Grammars

The `occ`s input file supports several % directives in addition to the ones discussed earlier. (All `occ`s directives are summarized in Table E.5. They will be described in this and subsequent sections.)

A definitions section for a small expression compiler is shown in Listing. E.8. The analyzer will recognize expressions made up of numbers, identifiers, parentheses, and addition and multiplication operators (+ and *). Addition is lower precedence than multiplication and both operators associate left to right. The compiler outputs code that evaluates the expression. The entire file, from which the definitions section was extracted, appears at the end of the `occ`s part of this appendix.

Terminal symbols are defined on lines one to six. Here, the `%term` is used for those terminal symbols that aren't used in expressions. `%left` is used to define left-associative operators, and `%right` is used for right-associative operators. (A `%nonassoc` directive is also supplied for declaring nonassociative operators.) The higher a `%left`, `%right`, or `%nonassoc` is in the input file, the lower its precedence. So, `PLUS` is lower precedence than `STAR`, and both are lower precedence than parentheses. A `%term` or `%token` is not needed if a symbol is used in a `%left`, `%right`, or `%nonassoc`. The precedence and associativity information is used by `occ`s to patch the parse tables created by an ambiguous input grammar so that the input grammar will be parsed correctly.

Specifying precedence
and associativity, `%term`,
`%left`, `%right`,
`%nonassoc`.

Table E.5. Occs % Directives and Comments

Directive	Description
%%	Delimits the three sections of the input file.
%{	Starts a code block. All lines that follow, up to a %} are written to the output file unchanged.
%}	Ends a code block.
%token	Defines a token.
%term	A synonym for %token.
/* */	C-like comments are recognized—and ignored—by occ s, even if they're outside of a %{ %} delimited code block.
%left	Specifies a left-associative operator.
%right	Specifies a right-associative operator.
%nonassoc	Specifies a nonassociative operator.
%prec	Use in rules section only. Modifies the precedence of an entire production to resolve a shift/reduce conflict caused by an ambiguous grammar.
%union	Used for typing the value stack.
%type	Attaches %union fields to nonterminals.

Listing E.8. *expr.y*—occs Definitions Section for a Small Expression Compiler

```

1  %term  ID      /* a string of lower-case characters */
2  %term  NUM     /* a number */
3
4  %left  PLUS    /* + */
5  %left  STAR    /* * */
6  %left  LP RP   /* ( ) */
7
8  %}
9  #include <stdio.h>
10 #include <ctype.h>
11 #include <stdlib.h>
12
13 extern char      *yytext;      /* In yylex(), holds lexeme */
14 extern char      *new_name();  /* declared at bottom of this file */
15
16 typedef char     *stype;      /* Value stack */
17 #define YYSTYPE   stype
18
19 #define YYMAXDEPTH 64
20 #define YYMAXERR   10
21 #define YYVERBOSE
22 %}

```

Most of the grammars in this book use recursion and the ordering of productions to get proper associativity and precedence. Operators handled by productions that occur earlier in the grammar are of lower precedence; left recursion gives left associativity; right recursion gives right associativity. This approach has its drawbacks, the main one being a proliferation of productions and a correspondingly larger (and slower) state machine. Also, in a grammar like the following one, productions like 2 and 4 have no purpose other than establishing precedence—these productions generate single-reduction states, described in Chapter Five.

0.	<i>s</i>	→	<i>e</i>
1.	<i>e</i>	→	<i>e + t</i>
2.			<i>t</i>
3.	<i>t</i>	→	<i>t * f</i>
4.			<i>f</i>
5.	<i>f</i>	→	(<i>e</i>)
6.			NUM
7.			ID

Occs lets you redefine the foregoing grammar as follows:

```
%term NUM ID

%left PLUS /* + */
%left STAR /* * */
%left LP RP /* ( ) */
%%

s : e ;
e : e PLUS e
| e STAR e
| LP e RP
| NUM
| ID
;
%%
```

The ambiguous grammar is both easier to read and smaller. It also generates smaller parse tables and a somewhat faster parser. Though the grammar is not LALR(1), parse tables can be built for it by using the disambiguation rules discussed in Chapter Five and below. If there are no ambiguities in a grammar, `%left` and `%right` need not be used—you can use `%term` or `%token` to declare all the terminals.

E.11.2 Attributes and the Occs Value Stack

Accessing bottom-up attributes, `$$`, `$1`, etc.

The `occ` parser automatically maintains a value stack for you as it parses. Moreover, it keeps track of the various offsets from the top of stack and provides a simple mechanism for accessing attributes. The mechanism is best illustrated with an example. In $S \rightarrow A B C$, the attributes can be accessed as follows:

```
S : A B C ;
$$ $1 $2 $3
```

That is, `$$` is the value that is inherited by the left-hand side after the reduce is performed. Attributes on the right-hand side are numbered from left to right, starting with the leftmost symbol on the right-hand side. The attributes can be used anywhere in a curly-brace-delimited code block. The parser provides a default action of `$$=$$1` which can be overridden by an explicit assignment in an action. (Put a `$$` to the left of an equals sign [just like a variable name] anywhere in the code part of a rule.)

The one exception to the `$$=$$1` rule is an ϵ production (one with an empty right-hand side)—there's no `$1` in an ϵ production. A reduce by an ϵ production pops nothing (because there's nothing on the right-hand side) and pushes the left-hand side. Rather than push garbage in this situation, `occ` duplicates the previous top-of-stack item in the push. `Yacc` pushes garbage.

Default action: `$$=$$1`.

Attributes and ϵ productions.

Typing the value stack.

By default, the value stack is of type `int`. Fortunately, it's easy to change this type. Just redefine `YYSTYPE` in the definitions section of the input file. For example, the expression compiler makes the value stack a stack of character pointers instead of `ints`.

with the following definitions:

```
%{
    typedef char      *stype;
    #define YYSTYPE stype
}%
```

The **typedef** isn't needed if the stack is redefined to a simple type (**int**, **long**, **float**, **double**, and so forth). Given this definition, you could use **\$1** to modify the pointer itself and ***\$1** to modify the object pointed to. (In this case, ***\$1** modifies the first character of the string.) A stack of structures could be defined the same way:

```
%{
    typedef struct
    {
        int      harpo;
        long     groucho;
        double   chico;
        char    zeppo[10];
    }
    stype;
#define YYSTYPE stype
}%
```

You can use **\$1.harpo**, **\$2.zeppo[3]**, and so forth, to access a field. You can also use the following syntax: **\$<harpo>\$**, **\$<chico>1**, and so forth—**\$<chico>1** is identical to **\$1.chico**. following statements:

Input	Expanded to
\$\$	Yy_val
\$N in the rules section	yysp[constant]
\$N in the code section	yysp[(Yy_rhslen - constant)]

yysp is the value-stack pointer. **Yy_rhslen** is the number of symbols on the right-hand side of the production being reduced. The *constant* is derived from *N*, by adjusting for the size of the right-hand side of the current production. For example, in a production like

```
S : A B C ;
```

\$1, **\$2**, and **\$3** evaluate to the following:

```
Yy_vsp[ -2 ] /* $1 */
Yy_vsp[ -1 ] /* $2 */
Yy_vsp[ -0 ] /* $3 */
```

Be careful of saying something like **\$\$=\$\$2** if you've defined **YYSTYPE** as a structure. The whole structure is copied in this situation. (Of course, if you want the whole structure to be copied...) Note that the default action (**\$\$=\$\$1**) is always performed before any of the actions are executed, and it affects the entire structure. Any specific action modifies the default action. This means that you can let some fields of a structure be inherited in the normal way and modify others explicitly. For example, using our earlier structure, the following:

```
{ $$ .harpo = 5; }
```

modifies the **harpo** field, but not the **groucho**, **chico**, or **zeppo** fields, which are inherited from **\$1** in the normal way.

Copying Structures.

The entire 24-byte structure is copied at every shift or reduce. Consequently, it's worth your effort to minimize the size of the value-stack elements. Note that the parser is built assuming that your compiler supports structure assignment. If this isn't the case, you'll have to modify the parser to use `memcpy()` to do assignments.

Occs' attribute support is extended from that of yacc in that dollar attributes can be used in both the rules and code sections of the input specification (`yacc` permits them only in the rules section). Occs treats `$0` as a synonym for `$$`. Finally, occs permits negative attributes. Consider the following:

```
s : X A B C
b : E           { $$ = $-1 + $-2 };
```

Negative attributes (`$-1`).

The `$-1` evaluates to the attributes for the symbol immediately below the `E` on the value stack. To see what the negative attributes are doing, consider the condition of the stack just before the reduce by $b \rightarrow E$:

`$1` references the attributes associated with `E` in the normal way; `$-1` references `A`'s attributes, and `$-2` references `X`'s attributes. Occs normally prints an error message if you try to reference an attribute that's outside the production (if you tried to use `$3` on a right-hand side that had only two elements, for example). It's possible to reference off the end of the stack if the numbers get too negative, however—no error message is printed in this case. For example, a `$-3` in the earlier example just silently evaluates to garbage (the start symbol doesn't have any attributes).

Now, look closely at the attribute-passing mechanism for the expression compiler, reproduced below:

```
e : e PLUS e { gen("%s += %s;", $1, $3); free_name($3); }
| e STAR e { gen("%s *= %s;", $1, $3); free_name($3); }
| LP e RP { $$ = $2;
}
| NUM       { gen("%s = %s;", $$ = new_name(), yytext); }
| ID        { gen("%s = %s;", $$ = new_name(), yytext); }
;
```

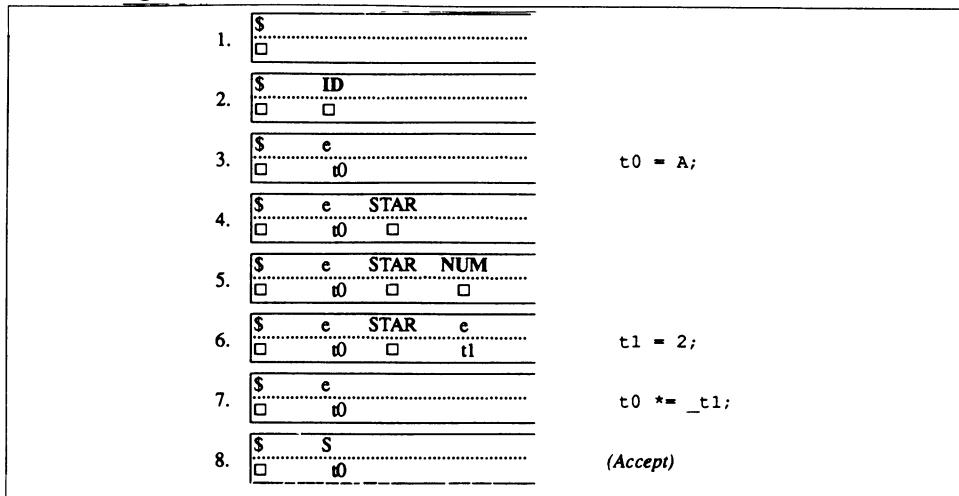
The $e \rightarrow \text{NUM}$ and $e \rightarrow \text{ID}$ actions on lines four and five are identical. They output a move instruction and then put the target of the move (the name of the rvalue) onto the value stack as an attribute. The `$$=$$2` in the $e \rightarrow \text{LP } e \text{ RP}$ action on the third line just moves the attribute from the e that's buried in the parentheses to the one on the left-hand side of the production. The top two productions are doing the real work. To see how they function, consider the sample parse of `A+2` in Figure E.5. The transition of interest is from stack picture six to seven. `t0` and `t1` were put onto the stack when the rvalues were created (and they're still there). So, we can find the temporaries used for the rvalues by looking at `$1` and `$2`. The inherited attribute is the name of the temporary that holds the result of the addition.

%union—Automatic field-name generation.

Because the value stack is more often than not typed as a union, a mechanism is provided to keep track of the various fields and the way that they attach to individual symbols. This mechanism assumes a common practice, that specific fields within a union are used only by specific symbols. For example, in a production like this:

```
e : e divop e
```

where `divop` could be either `/` or `%` (as in C), you would need to store one of two attributes on the value stack. The attribute attached to the `e` would be a pointer to a string holding the name of a temporary variable; the attribute attached to the `divop` would be `'/'` for divide and `'%'` for a modulus operation. To do the foregoing, the value stack must be typed to be a union of two fields, one a character pointer and the other an `int`. You could do this using an explicit redefinition of `YYSTYPE`, as discussed earlier:

Figure E.5. A Parse of A*2

```
typedef union
{
    char *var_name ;
    int op_type ;
}
yystype;
```

```
#define YYSTYPE yystype
```

but the **%union** directive is a better choice. Do it as follows:

```
%union {
    char *var_name ;
    int op_type ;
}
```

A **%union** activates an automatic field-name generation feature that is not available if you just redefine **YYSTYPE**. Given the earlier production (**e** : **e** DIVOP **e**), we'd like to be able to say **\$1**, **\$2**, and so forth, without having to remember the names of the fields. That is, you want to say

```
$1 = new_name();
```

rather than

```
$1.var_name = new_name();
```

Do this, first by using the **%union**, and then attaching the field names to the individual terminal and nonterminal symbols with a **<name>** operator in token-definition directive (**%term**, **%left**, etc.) or a **%type** directive. The **name** can be any field name in the **%union**. For example:

```
%term <op_type> DIVOP
%type <var_name> e
```

attaches the **op_type** field to all **DIVOPs** and the **var_name** field to all **e's**. That is, if a **<name>** is found in a **%term**, **%token**, **%left**, **%right**, or **%nonassoc** directive, the indicated field name is automatically attached to the specified token. The angle brackets are part of the directive. They must be present, and there may not be any white space between them and the actual field name. The **%type** directive is used to attach a field name to a nonterminal.

<name> and **%type**

Listing E.9. *union.y*— Using %union: An Example

```

1  %union {
2      int   op_type;
3      char  *var_name;
4  }
5
6  %term DIVOP                                /* / or %
7  %type <var_name> e statement
8  %type <op_type> divop
9
10 %%
11 goal      : statement
12 ;
13
14 statement : e divop e
15 {
16     if( $2 == '/' )
17         gen( "%s /= %s;", $1, $3 )
18     else
19         gen( "%s != %s;", $1, $3 )
20
21     free_name( $3 );
22     $$ = $1
23 }
24 ;
25 divop    : DIVOP  ( $$ = yytext; )
26 ;
27 %%

```

A real example of all the foregoing is in Listing E.9. The grammar recognizes a single statement made up of two numbers separated by a DIVOP, which can be a / or %. The \$\$ in the divop rule on line 25 references the op_type field of the union because divop is attached to that field on line eight. Similarly, the \$2 on line 16 references the op_type field because \$2 corresponds to a divop in this production, and divop was attached to op_type on line eight. The \$1 and \$3 on lines 17, 19, 21, and 22 reference the var_name field because they correspond to an e in the production and were attached to e on line seven. Similarly the \$\$ on line 22 references var_name because it corresponds to a statement, which was also attached to a var_name on line seven.

The field stays active from the point where it's declared to the end of the current directive. For example, in

```
%token ONE <field1> TWO THREE <field2> FOUR
```

<field1> is attached to TWO and THREE, and <field2> is attached to FOUR. ONE has no field attached to it. If a \$\$ or \$N corresponds to a symbol that has no explicit %union field attached to it, a default int-sized field is used and a warning message is printed.¹² Also, note that the nonterminal name on the left-hand side of a production must be found in a previous %type in order to use \$\$ in the corresponding right-hand side. The %type directive can have more than one nonterminal symbol listed in it, but it may have only one <name> and the name must be the first thing in the list. Similarly, the <name> in a %left and so forth applies only to those terminals that follow it in the list.¹³

\$<field>N

The \$<field>N syntax described earlier can also be used to access a field of a %union. This syntax takes precedence over the default mechanism, so you can use it

12. Yacc doesn't print the warning.

13. Yacc is very picky about type matching when you use this mechanism. Since this pickiness is now supported by most compilers, I've not put any type checking into occs. It just stupidly makes the substitutions discussed earlier.

when you want to override the default field for a particular symbol. It's also useful in imbedded actions where there's no mechanism for supplying a field name to `$$`. For example, in:

```
%union { integer }
...
%%
x : a { $<integer>$ = 1; } b { foo( $<integer>2 ); }
```

The `$<integer>$` is used to put 1 into the `integer` field of the synthesized attribute. That number is later referenced using the `$<integer>2` in the second action. Since the field type is an `integer`, you could also use the default field as follows:

```
x : a { $$ = 1; } b { foo($2); }
```

but `occ`s prints a warning message when no field is specified, as is the case here. (`Yacc` doesn't print the warning).

`Occs` translates all `%unions` into a `typedef` and redefinition of `YYSTYPE`. The foregoing example is output as follows:

`typedef` for `%union`:
`yystype`.

```
typedef union
{
    int    yy_def;
    char   *var_name ;
    int    op_type ;
}
yystype;

#define YYSTYPE yystype
```

The `YYSTYPE` definition is generated automatically. Both it and the `yystype` type name will be useful later on. The `yy_def` field is used for the default situation, where the symbol specified in a dollar attribute has not been attached to an explicit symbol. For example, if the following production is found without a preceding `%type <field>` etc:

```
t : e PLUS e { $$ = $1 + $3; } ;
```

the `yy_def` field will be used for both `$$` (because there's no field attached to the `t`) and for `$1` and `$2` (because there's no field attached to the `e` either).

E.11.3 Printing the Value Stack

If you're using the default stack type of `int` you don't have to do anything special to print the value stack in the debugging environment because the default stack-printing routine, `yypstk()`, will be called in from the library. The library version of this routine assumes a stack of integers.

Printing an `int` value stack.

If you've changed the stack type, either by redefining the `YYSTYPE` or macro or by using a `%union`, you'll need to provide a version of the `yypstk()` function for the debugger to use. The function is passed a pointer to a value-stack element and a second pointer to the equivalent symbol-stack element. It must return a pointer to a string that represents the value-stack element.

Printing a non-`int` value stack.

If the value stack is a simple type, such as a stack of character pointers, then printing the stack contents is easy:

```

typedef char *stack_type ;
#define YYSTYPE stack_type ;
%%
...
/* productions go here */
%%
char *yypstk( value, symbol )
YYSTYPE *value; /* pointer to a character pointer */
char *symbol;
{
    return *value;
}

```

When the value stack is typed with a %union, things are a little more complicated. Listing E.10 shows how the value stack for the earlier example would be printed. The %union forces the value stack to be an array of YYSTYPEs, and a pointer to one element of this array is passed to yypstk(). The second, symbol argument to yypstk() is used to select the correct field of the union.

Listing E.10. *yypstk2.c*—Printing Attributes for a %union

```

1  %union {
2      int    op_type;
3      char   *var_name;
4  }
5
6  %term DIVOP                                /* / or % */
7  %type <var_name> e statement
8  %type <op_type>  divop
9  %%
10
11 %%
12 yypstk( value, symbol )
13 YYSTYPE *value;
14 char   *symbol;
15 {
16     static char buf[80];
17
18     if ( !strcmp(symbol,"e") || !strcmp(symbol,"statement") )
19         return value->var_name ;
20
21     if ( !strcmp(symbol,"divop") )
22     {
23         sprintf( buf, "%c", value->op_type );
24         return buf;
25     }
26     else
27         return "---"; /* other symbols don't have attributes in this */
28                                         /* application. */
29 }

```

If you were using the default, **int** attribute field in the union, the **return** "---" statement on line 27 would be replaced with the following:

```

sprintf( buf, "%d", value->yy_def );
return buf;

```

E.11.4 Grammatical Transformations

Imbedded actions (ones that are in the middle of a production) can be used in an **occ**s grammar, but you have to be careful with them. The problem is that actions can be executed only during a reduction. Consequently, if you put one in the middle of a production, **occ**s has to shuffle the grammar by adding an ϵ production. For example:

```
s : a { action(1); } b c { action(2); }
;
```

is modified by **occ**s as follows:

```
s      : a 0001 b c { action(2); } ;
0001 : { action(1); }
;
```

Unfortunately, that extra ϵ production can introduce shift/reduce conflicts into the state machine. It's best, if possible, to rewrite your grammar so that the imbedded production isn't necessary:

```
s      : prefix b c { action(2); } ;
prefix : a           { action(1); }
;
```

Using the attribute mechanism is also a little tricky if you imbed an action in a production. Even though imbedded actions are put in their own production, the $\$1$, $\$2$, and so forth reference the parent production. That is, in:

```
s : a b { $$ = $1 + $2 } c { $$ = $1 + $2 + $3 + $4 } ;
;
```

the $\$1$ and $\$2$ in both actions access the attributes associated with *a* and *b*. $\$\$$ in the left action is accessed by $\$3$ in the right action. (That is, this $\$\$$ is actually referencing the 0001 left-hand side inserted by **occ**s.) The $\$\$$ in the right action is attached to *s* in the normal way. $\$4$ accesses *c*. Note that $\$3$ and $\$4$ are illegal in the left action because they won't be on the parse stack when the action is performed. An error message is generated in this case, but only if the reference is in the actual grammar section of the **yacc** specification. Illegal stack references are silently accepted in the final code section of the input file.

Occs supports the two non-**yacc** transformations. Brackets are used to designate optional parts of a production. The following input:

```
s : a [ b ] c ;
;
```

is translated internally to:

```
s      : a 001 c
;
001   : b
| /* epsilon */
;
```

Optional subexpressions,
[...].

Note that attributes in optional productions can be handled in unexpected ways (which actually make sense if you consider the translation involved). For example, in:

```
s -> a [b c { $$ = $1 + $2; }] d { $$ = $1 + $2 + $3; }
```

The $\$\$=\$1+\2 in the optional production adds the attributes associated with *b* and *c* and attaches the result to the entire optional production. The action on the right adds together the attributes associated with *a*, the entire optional production *bc* and *d*. That is, $\$2$ in the right production is picking up the $\$\$$ from the optional production. Note that the $\$2$ used in the right action is garbage if the ϵ production was taken in the optional part of the production. As a consequence, optional productions are typically not used when attributes need to be passed.

Optional productions nest. The following is legal:

```
s -> a [ b [c] [d [e]] ] f
```

though almost incomprehensible—I wouldn't recommend using it. The maximum nesting level is 8. Note that optional subexpressions can introduce duplicate productions. That is:

```
s : b [c] d
| e [c]
;
```

creates:

```
s : b 001 d
| e 002
;
001 : c
| /* epsilon */
;
002 : c
| /* epsilon */
;
```

It's better to use the following in the original grammar:

```
s : b opt_c d
| e opt_c
;
opt_c : c
| /* empty */
;
```

Also note that

```
s : [x] ;
```

is acceptable but needlessly introduces an extra production:

```
s : 001 ;
001 : x | /* empty */ ;
```

It's better to use:

```
s : x | /* empty */ ;
```

Repeating subexpressions. [...]*.

Adding a star after the right bracket causes the enclosed expression to repeat zero or more times. A left-associative list is used by **occ**; a right-associative list is used by **LLama**. The internal mappings for all kinds of brackets are shown in the Table E.6.

You can't use this mechanism if you need to pass attributes from *b* back up to the parent, primarily because you can't attach an action to the added ε production. The extra ε production may also add shift/reduce or reduce/reduce conflicts to the grammar. Be careful.

Some examples—a comma-separated list that has at least one element is:

```
s : a [ COMMA a ]* ;
```

A dot-separated list with either one or two elements is:

```
s : a [DOT a] ;
```

One or more repetitions of *b* is:

```
s : a b [b]* c ;
```

Table E.6. Occs Grammatical Transformations

Input	Output
s : a [b] c ;	s : a 001 c 001 : b /* epsilon */ ;
s : a [b]* c ;	s : a 001 c 001 : 001 b /* epsilon */ ; <i>(occ version)</i>
s : a [b]* c ;	s : a 001 c 001 : b 001 /* epsilon */ ; <i>(llama version)</i>

E.11.5 The *yyout.sym* File

You can see the transformations made to the grammar, both by adding imbedded actions and by using the bracket notation for optional and repeating sub-productions, by looking in the symbol-table file, *yyout.sym*, which will show the transformed grammar, not the input grammar. The symbol-table file is generated if *-D*, *-S*, or *-s* is specified on the command line. A *yyout.sym* for the grammar in Listing E.2 on page 841 is in Listing E.11. (It was generated with *occ -S expr.y*.)

Generating the symbol table, *yyout.sym*, *-D*, *-S*, *-s*.

Listing E.11. *yyout.sym*—Occs Symbol-Table Output

```

1 ----- Symbol table -----
2
3 NONTERMINAL SYMBOLS:
4
5 e (257)    <>
6     FIRST : ID NUM LP
7         5: e -> ID
8         4: e -> NUM
9         3: e -> LP e RP .....PREC 3
10        2: e -> e STAR e .....PREC 2
11        1: e -> e PLUS e .....PREC 1
12
13 s (256)  (goal symbol) <>
14     FIRST : ID NUM LP
15         0: s -> e
16
17 TERMINAL SYMBOLS:
18
19 name          value  prec  assoc   field
20 STAR           4      2      1      <>
21 RP              6      3      1      <>
22 PLUS            3      1      1      <>
23 NUM             2      0      -      <>
24 ID              1      0      -      <>
25 LP              5      3      1      <>

```

Taking a few lines as a representative sample:

```

e (257)    <>
FIRST : ID NUM LP
5: e -> ID
4: e -> NUM
3: e -> LP e RP .....PREC 3

```

A %union<field> in
yyout.sym.

The FIRST set in
yyout.sym.

Right-hand sides in
yyout.sym, production
numbers.

Precedence, associativity
in *yyout.sym*.

Modified productions are
shown.

the top line is the name and internal, tokenized value for the left-hand side. The <> on the right contains the field name assigned with a %type directive. Since there was none in this example, the field is empty. Had the following appeared in the input:

```
%type <field> e
```

the <field> would be on the corresponding line in the symbol table. The next line is the symbol's FIRST set. (The list of terminal symbols that can appear at the far left of a parse tree derived from the nonterminal.) That is, it's the set of tokens which can legitimately be next in the input when we're looking for a given symbol in the grammar. Here, we're looking for an *e*, and an ID, NUM, or LP can all start an *e*. The FIRST sets are output only if the symbol table is created with -S.

The next few lines are the right-hand sides for all productions that share the single left-hand side. The number to the left of the colon is the production number. These are assigned in the order that the productions were declared in the source file. The first production is 0, the second is 1, and so forth. Note that the production that has the goal symbol on its left-hand side is always Production 0. Production numbers are useful for setting breakpoints in the interactive debugging environment. (You can set a breakpoint on reduction by a specific production, provided that you know its number.) The PREC field gives the relative precedence of the entire production. This level is used for resolving shift/reduce conflicts, also discussed below. Note that the productions shown in this table are the ones actually used by occs—any transformations caused by imbedded actions or the [] operator are reflected in the table. Productions are sorted alphabetically by left-hand side.

The second part of the table gives information about the terminal symbols. For example:

name	value	prec	assoc	field
PLUS	3	1	1	<>
NUM	2	0	-	<>

Here, 3 is the internal value used to represent the PLUS token. It is also the value assigned to a PLUS in *yyout.h*. The prec field gives the relative precedence of this symbol, as assigned with a %left, %right, or %nonassoc directive. The higher the number, the higher the precedence. Similarly, the assoc is the associativity. It will have one of the following values:

l	left associative
r	right associative
n	nonassociative
-	Associativity is not specified. The token was declared with a %term or %token rather than a %left, %right, or %nonassoc directive.

Finally, the field column lists any %union fields assigned with a <name> in the %left, %right, and %nonassoc directives. A <> is printed if there is no such assignment.

E.11.6 The *yyout.doc* File

The *yyout.doc* file holds a description of the LALR(1) state machine used by the occs-generated parser. It is created if -v or -V is present on the command line. The *yyout.doc* file for the grammar in Listing E.2 on page 841 is shown in Table E.7. Figure E.6 shows the state machine in graphic form.

This machine has ten states, each with a unique number (running from 0 to 9). The top few lines in each state represent the LALR(1) kernel items. You can use them to see

LALR(1) states in
yyout.doc

the condition of parse stack when the current state is reached. For example, the header from State 9 looks like this:

```
State 9:
e->e .PLUS e
e->e PLUS e .  [$ PLUS STAR RP]
e->e .STAR e
```

The dot is used to mark the current top of stack, so everything to the left of the dot will be on the stack. The top-of-stack item in State 9 is an *e* because there's an *e* immediately to the left of the dot. The middle line is telling us that there may also be a PLUS and another *e* under the *e* at top of stack.

If the dot is at the far right (as it is on the middle line), then a handle is on the stack and the parser will want to reduce. The symbols in brackets to the right of the production are a list of those symbols that (at least potentially) cause a reduction if they're the next input symbol and we are in State 9. (This list is the LALR(1) lookahead set for the indicated LR item, as discussed in Chapter Five.) The lookahead set is printed for every kernel item (as compared to only those items with the dot at the far right) if $-V$ is used rather than $-v$. Note that these are just potential reductions, we won't necessarily do the reduction on every symbol in the list if there's a conflict between the reduction and a potential shift. A $\$$ is used in the lookahead list to represent the end-of-input marker.

The next lines show the possible transitions that can be made from the current state. There are four possibilities, which will look something like the following:

Reduce by 3 on PLUS

says that a reduction by Production 3 occurs if the next input symbol is a PLUS.

Shift to 7 on STAR

says that a 7 is shifted onto the stack (and the input advanced) if the next input symbol is a STAR.

Goto 4 on e

takes care of the push part of a reduction. For example, starting in State 0, a NUM in the input causes a shift to State 2, so a 2 is pushed onto the stack. In State 2, a PLUS in the input causes a reduce by Production 4 ($e \rightarrow \text{NUM}$), which does two things. First, the 2 is popped, returning us to State 0. Next, the parser looks for a goto transition (in State 0) that is associated with the left-hand side of the production by which we just reduced. In this case, the left-hand side is an ϵ , and the parser finds a Goto 4 on e in State 0, so a 4 is pushed onto the stack as the push part of the reduction. The final possibility,

Accept on end of input

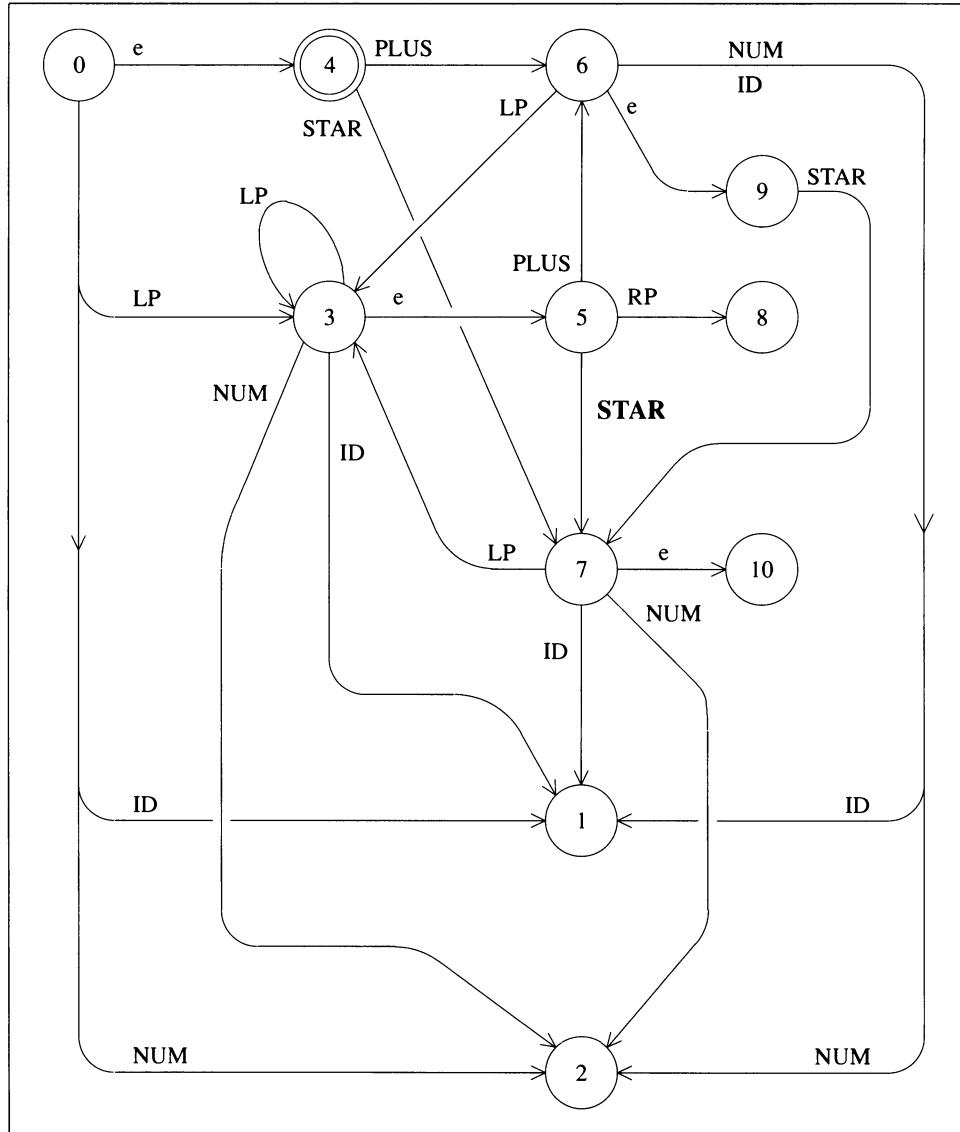
says that if the end-of-input marker is found in this state, the parse is successfully terminated.

Symbols that cause reductions in *yyout.doc* state.

Transitions in *yyout.doc*, shift, reduce, accept.

Table E.7. yyout.doc (Generated from Listing E.2)

<p>State 0: $s \rightarrow e$ Shift to 1 on ID Shift to 2 on NUM Shift to 3 on LP Goto 4 on e</p>	<p>State 6: $e \rightarrow e \text{ PLUS } .e$ Shift to 1 on ID Shift to 2 on NUM Shift to 3 on LP Goto 9 on e</p>
<p>State 1: $e \rightarrow ID \ . \ [\\$ \text{ PLUS STAR RP }]$ Reduce by 5 on End of Input Reduce by 5 on PLUS Reduce by 5 on STAR Reduce by 5 on RP</p>	<p>State 7: $e \rightarrow e \text{ STAR } .e$ Shift to 1 on ID Shift to 2 on NUM Shift to 3 on LP Goto 10 on e</p>
<p>State 2: $e \rightarrow NUM \ . \ [\\$ \text{ PLUS STAR RP }]$ Reduce by 4 on End of Input Reduce by 4 on PLUS Reduce by 4 on STAR Reduce by 4 on RP</p>	<p>State 8: $e \rightarrow LP \ e \ RP \ . \ [\\$ \text{ PLUS STAR RP }]$ Reduce by 3 on End of Input Reduce by 3 on PLUS Reduce by 3 on STAR Reduce by 3 on RP</p>
<p>State 3: $e \rightarrow LP \ .e \ RP$ Shift to 1 on ID Shift to 2 on NUM Shift to 3 on LP Goto 5 on e</p>	<p>State 9: $e \rightarrow e \ .PLUS \ e$ $e \rightarrow e \text{ PLUS } e \ . \ [\\$ \text{ PLUS STAR RP }]$ $e \rightarrow e \ .STAR \ e$ Reduce by 1 on End of Input Reduce by 1 on PLUS Shift to 7 on STAR Reduce by 1 on RP</p>
<p>State 4: $s \rightarrow e \ . \ [\\$ \]$ $e \rightarrow e \ .PLUS \ e$ $e \rightarrow e \ .STAR \ e$ Accept on end of input Shift to 6 on PLUS Shift to 7 on STAR</p>	<p>State 10: $e \rightarrow e \ .PLUS \ e$ $e \rightarrow e \ .STAR \ e$ $e \rightarrow e \text{ STAR } e \ . \ [\\$ \text{ PLUS STAR RP }]$ Reduce by 2 on End of Input Reduce by 2 on PLUS Reduce by 2 on STAR Reduce by 2 on RP</p>
<p>State 5: $e \rightarrow e \ .PLUS \ e$ $e \rightarrow e \ .STAR \ e$ $e \rightarrow LP \ e \ .RP$ Shift to 6 on PLUS Shift to 7 on STAR Shift to 8 on RP</p>	<p>6/254 terminals 2/256 nonterminals 6/512 productions 11 states</p>

Figure E.6. State Machine Represented by the *yyout.doc* File in Listing E.11

E.11.7 Shift/Reduce and Reduce/Reduce Conflicts

One of the main uses of the *yyout.doc* file is to see how shift/reduce and reduce/reduce conflicts are solved by *occs*. You should never let a WARNING about an inadequate state go by without looking in *yyout.doc* to see what's really going on.

Ococs uses the disambiguating rules discussed in Chapter Five to resolve conflicts. Reduce/reduce conflicts are always resolved in favor of the production that occurred earlier in the grammar. Shift/reduce conflicts are resolved as follows:

Resolving shift/reduce and reduce/reduce conflicts.

- (1) Precedence and associativity information is assigned to all terminal symbols using `%left`, `%right`, and `%nonassoc` directives in the definitions part of the input file. The directives might look like this:

```
%left PLUS MINUS
%left TIMES DIVIDE
%right ASSIGN
```

The higher a directive is in the list, the lower the precedence. If no precedence or associativity is assigned, a terminal symbol will have a precedence level of zero (very low) and be nonassociative.

Using `%prec`.

Productions are assigned the same precedence level as the rightmost terminal symbol in the production. You can override this default with a `%prec TOKEN` directive to the right of the production. (It must be between the rightmost symbol in the production and the semicolon or vertical bar that terminates the production). TOKEN is a terminal symbol that was declared with a previous `%left`, `%right`, or `%nonassoc`, and the production is assigned the same precedence level as the indicated token. `Occs`, but not `yacc`, also allows statements of the form

```
%prec number
```

where the *number* is the desired precedence level (the higher the number, the higher the precedence). The *number* should be greater than zero.

- (2) When a shift/reduce conflict is encountered, the precedence of the terminal symbol to be shifted is compared with the precedence of the production by which you want to reduce. If the terminal or production is of precedence zero, then resolve in favor of the shift.
- (3) Otherwise, if the precedences are equal, resolve using the following table:

<i>associativity of lookahead symbol</i>	<i>resolve in favor of</i>
left	reduce
right	shift
nonassociative	shift

- (4) Otherwise, if the precedences are not equal, use the following table:

<i>precedence</i>	<i>resolve in favor of</i>
lookahead symbol < production	reduce
lookahead symbol > production	shift

The `%prec` directive can be used both to assign a precedence level to productions that don't contain any terminals, and to modify the precedences of productions in which the rightmost nonterminal isn't what we want. A good example is the unary minus operator, used in the following grammar:

```
%term      NUM
%left     MINUS PLUS
%left     TIMES
%nonassoc VERY_HIGH
%%
s : e ;
e : e PLUS e
| e MINUS e
| e TIMES e
| MINUS e  %prec VERY_HIGH
| NUM
;
%%
```

Here, the `%prec` is used to force unary minus to be higher precedence than both binary

minus and multiplication. `VERY_HIGH` is declared only to get another precedence level for this purpose. `Occs` also lets you assign a precedence level directly. For example,

```
| MINUS e      %prec 4
```

could have been used in the previous example.¹⁴ C's `sizeof` operator provides another example of how to use `%prec`. The precedence of `sizeof` must be defined as follows:

```
expression : SIZEOF LP type_name RP      %prec SIZEOF
;
```

in order to avoid incorrectly assigning a `sizeof` the same precedence level as a right parenthesis.

The precedence level and `%prec` operator can also be used to resolve shift/reduce conflicts in a grammar. The first technique puts tokens other than operators into the precedence table. Consider the following state (taken from `yyout.doc`).

Using `%prec` to resolve shift/reduce conflicts.

```
WARNING: State 5: shift/reduce conflict ELSE/40 (choose shift)
```

```
State 5:
stmt-> IF LP expr RP stmt.
| IF LP expr RP stmt. ELSE stmt
```

The default resolution (in favor of the shift) is correct, but it's a good idea to eliminate the warning message (because you don't want to clutter up the screen with harmless warning messages that will obscure the presence of real ones). You can resolve the shift/reduce conflict as follows. The conflict exists because `ELSE`, not being an operator, is probably declared with a `%term` rather than a `%left` or `%right`. Consequently, it has no precedence level. The precedence of the first production is taken from the `RP` (the rightmost terminal in the production), so to resolve in favor of the shift, all you need do is assign a precedence level to `ELSE`, making it higher than `RP`. Do it like this:

```
%left      LP RP /* existing precedence of LP */           */
%nonassoc ELSE /* higher precedence because it follows LP */
```

Though you don't want to do it in the current situation, you *could* resolve in favor of the reduce by reversing the two precedence levels. (Make `ELSE` lower precedence than `LP`).

The second common situation is illustrated by the following simplification of the C grammar used in Chapter Six. (I've left out some of the operators, but you get the idea.)

```
function_argument
: expr
| function_argument COMMA expr /* comma separates arguments */
;
expr
: expr STAR      expr
| expr COMMA     expr          /* comma operator */
| expr DIVOP    expr
...
| term
;
```

A shift/reduce conflict is created here because of the `COMMA` operator (the parser

¹⁴ Yacc doesn't permit this. You have to use a bogus token.

doesn't know if a comma is an operator or a list-element separator), and this conflict is displayed in *yyout.doc* as follows:

```
WARNING: State 170: shift/reduce conflict COMMA/102 (choose shift)
State 170:
    function_argument-> expr. (prod. 102, prec. 0)
        [COMMA RP ]
    expr-> expr. STAR expr
    expr-> expr. COMMA expr
    expr-> expr. DIVOP expr
```

The problem can be solved by assigning a precedence level to production 102: *function_argument* → *expr*. (It doesn't have one because there are no terminal symbols in it). You can resolve in favor of the reduce (the correct decision, here) by giving the production a precedence level greater than or equal to that of the comma. Do this in the input file as follows:

```
expr_list
    : expr                                %prec COMMA
    | expr_list COMMA expr
    ;
```

Similarly, you could resolve in favor of the shift by making the production lower precedence than the COMMA (by replacing the COMMA in the %prec with the name of a lower-precedence operator). Since the comma is the lowest-precedence operator in C, you'd have to do it here by creating a bogus token that has an even lower precedence, like this:

```
%nonassoc VERY_LOW /* bogus token (not used in the grammar) */
%left      COMMA
...
%%
...
expr_list
    : expr                                %prec VERY_LOW
    | expr_list COMMA expr
    ;
```

Shift/reduce and reduce/reduce conflicts are often caused by the implicit ε productions that are created by actions imbedded in the middle of a production (rather than at the far right), and the previous techniques can not be used to resolve these conflicts because there is no explicit production to which a precedence level can be assigned. For this reason, it's best to use explicit ε productions rather than imbedded actions. Translate:

```
x : a {action();} b ;
```

to this:

```
x      : a action b ;
action : /* empty */ {action();} ;
```

or to this:

```
x      : a' b
a'     : a {action();}
```

These translations probably won't eliminate the conflict, but you can now use %prec to resolve the conflict explicitly.

It's not always possible to do a translation like the foregoing, because the action may have to access the attributes of symbols to its left in the parent production. You can sometimes eliminate the conflict just by changing the position of the action in the production, however. For example, actions that follow tokens are less likely to introduce

conflicts than actions that precede them. Taking an example from the C grammar used in Chapter Six, the following production generated 40-odd conflicts:

```
and_list
: and_list {and($1);} ANDAND binary { and($4); $$=NULL; }
| binary
;
```

but this variant generated no conflicts and does the same thing:

```
and_list
: and_list ANDAND {and($1);} binary { and($4); $$=NULL; }
| binary
;
```

E.11.8 Error Recovery

One of the ways that **occ**s differs considerably from **yacc** is the error-recovery mechanism.¹⁵ **Occs** parsers do error recovery automatically, without you having to do anything special. The *panic-mode* recovery technique that was discussed in Chapter 5 is used. It works as follows:

- (0) An error is triggered when an error transition is read from the parse table entry for the current input and top-of-stack symbols (that is, when there's no legal outgoing transition from a state on the current input symbol).
- (1) Discard the state at the top of stack.
- (2) Look in the parse table and see if there's a legal transition on the current input symbol and the uncovered stack item. If so, we've recovered and the parse is allowed to progress, using the modified stack. If there's no legal transition, and the stack is not empty, go to 1.
- (3) If all items on the stack are discarded, restore the stack to its original condition, discard an input symbol, and go to 1.

The algorithm continues either until it can start parsing again or the entire input file is absorbed. In order to avoid cascading error messages, messages are suppressed if a second error happens right on the tail of the first one. To be more exact, no messages are printed if an error happens within five parse cycles (five shift or reduce operations) of the previous error. The number of parse cycles can be changed with a

```
%{ #define YYCASCADE desired_value %}
```

in the specifications section. Note that errors that are ignored because they happen too soon aren't counted against the total defined by YYMAXERR.

Panic-mode error recovery.

Avoiding cascading error messages, YYCASCADE.

E.11.9 Putting the Parser and Actions in Different Files

Unfortunately, **occ**s can take a long time to generate the parse tables required for a largish grammar. (Though it usually takes less time for **occ**s to generate *yyout.c* than it does for Microsoft C to compile it.) To make program development a little easier, a mechanism is provided to separate the table-making functions from the code-generation

¹⁵. **Yacc** uses a special `error` token, which the parser shifts onto the stack when an error occurs. You must provide special error-recovery productions that have `error` tokens on the right-hand sides. The mechanism is notoriously inadequate, but it's the only one available if you're using **yacc**. See [Schreiner] pp. 65–82 for more information.

Using the `-p` and `-a` command-line switches.

functions. The `-p` command-line switch causes `occ`s to output the tables and parser only (in a file called `yyout.c`). Actions that are part of a rule are not output, and the third part of the `occ`s input file is ignored. When `-a` is specified, only the actions are processed, and tables are not generated. A file called `yyact.c` is created in this case.

Once the grammar is stable you can run `occ`s once with `-p` to create the tables. Thereafter, you can run the same file through `occ`s with `-a` to get the actions. You now have two files, `yyout.c` and `yyact.c`. Compile these separately, and then link them together. If you change the actions (but not the grammar), you can recreate `yyact.c` using `occ`s `-a` without having to remake the tables. Remember that actions that are imbedded in the middle of a production will effectively modify the grammar. If you modify the position of an action in the grammar, you'll have to remake the tables (but not if you modify only the action code). On the other hand, actions added to or removed from the far right of a production will not affect the tables at all, so can be modified, removed, or added without needing to remake the tables.

`YYACTION, YYPARSER.`

`Yy_val`, `Yy_vsp`, and
`Yy_rhslen` made public
by `-p` or `-a`.

The first, definitions part of the `occ`s input file is always output, regardless of the presence of `-a` or `-p`. The special macro `YYPARSER` is generated if a parser is present in the current file, `YYACTION` is generated if the actions are present. (Both are defined when neither switch is specified.) You can use these in conjunction with `#ifdef`s in the definitions section to control the declaration of variables, and so forth (to avoid duplicate declarations). It's particularly important to define `YYSTYPE`, or put a `%union`, in both files—if you're not using the default `int` type, that is—otherwise, attributes won't be accessed correctly. Also note that three global variables whose scope is normally limited to `yyout.c`—`Yy_val`, `Yy_vsp`, and `Yy_rhslen`—are made public if either switch is present. They hold the value of `$$`, the value-stack pointer, and the right-hand side length.

Listing E.12 shows the definitions section of an input file that's designed to be split up in this way.

Listing E.12. Definitions Section for `-a/-p`

```

1  %{
2  #ifdef YYPARSER           /* If parser is present, declare variables.      */
3  #define CLASS
4  #define _(x) x
5  #else                      /* If parser is not present, make them externs. */
6  #define CLASS extern
7  #define _(x)
8  #endif
9
10 CLASS int x _( = 5 );    /* Evaluates to "int x = 5;" in yyparse.c and to */
11                                /* "extern int x;" in yyacts.c.                      */
12 %
13
14 %union {                  /* Either a %union or a redefinition of YYSTYPE */
15     ...                   /* should go here.                                */
16 }
17
18 %term ...                 /* Token definitions go here.                      */
19 %left ...
20 %%

```

E.11.10 Shifting a Token's Attributes

It is sometimes necessary to attach an attribute to a token that has been shifted onto the value stack. For example, when the lexical analyzer reads a **NUMBER** token, it would be nice to shift the numeric value of that number onto the value stack. One way to do this is demonstrated by the small desk-calculator program shown in the the **occs** input file in Listing E.13. The L^EX lexical analyzer is in Listing E.14. This program sums together a series of numbers separated by plus signs and prints the result. The input numbers are converted from strings to binary by the `atoi()` call on line 12 and the numeric value is also pushed onto the stack here. The numeric values are summed on line eight, and the result is printed on line four. The difficulty, here, is that you need to introduce an extra production on line 12 so that you can shift the value associated with the input number, making the parser both larger and slower as a consequence.

Using an extra production
to shift attribute for a
token.

Listing E.13. Putting Token Attributes onto the Value Stack—Using A Reduction

```

1 %term NUMBER      /* a collection of one or more digits    */
2 %term PLUS        /* a + sign                                */
3 %%
4 statement : expr statement      { printf("The sum is %d\n", $1); }
5           | /* empty */
6           ;
7
8 expr     : expr PLUS number    { $$ = $1 + $3;          }
9           | /* empty */
10          ;
11
12 number   : NUMBER            { $$ = atoi(yytext);      }
13          ;
14 %%

```

Listing E.14. Lexical Analyzer for Listing E.13

```

1 #include "yyout.h"      /* token definitions */
2 %%
3 [0-9]+  return NUMBER;
4 \+      return PLUS;
5 .       ;               /* ignore everything else */
6 %%

```

An alternate solution to the problem is possible by making the lexical analyzer and parser work together in a more integrated fashion. The default action in a shift (as controlled by the YYSHIFTACT macro described earlier) is to push the contents of a variable called `yylval` onto the value stack. The lexical analyzer can use this variable to cause a token's attribute to be pushed as part of the shift action (provided that you haven't modified YYSHIFTACT anywhere).¹⁶ The procedure is demonstrated in Listings E.15 and E.16, which show an alternate version of the desk calculator. Here, the lexical analyzer assigns the numeric value (the attribute of the **NUMBER** token) to `yylval` before

Using `yylval` to pass
attributes from scanner
to parser.

16. Note that `yacc` (but not **occs**) also uses `yylval` to hold `$$`, so it should never be modified in a `yacc` application, except as described below (or you'll mess up the value stack).

returning the token. That value is shifted onto the value stack when the token is shifted, and is available to the parser on line eight.

Listing E.15. Passing Attributes from the Lexical Analyzer to the Parser

```

1 %term NUMBER      /* a collection of one or more digits    */
2 %term PLUS        /* a + sign                                */
3 %%
4 statement : expr statement      { printf("The sum is %d\n", $1); }
5           | /* empty */
6           ;
7
8 expr      : expr PLUS NUMBER     { $$ = $1 + $3; }
9           | /* empty */
10          ;
11 %%

```

Listing E.16. Lexical Analyzer for Listing E.15

```

1  %{
2  #include "yyout.h"      /* token definitions                      */
3  extern int yylval;      /* declared by parser in occs output file */
4  %}
5  %%
6 [0-9]+ yylval = atoi(yytext); /* numeric value is synthesized in shift */
7  return NUMBER;
8
9 \+
10 .
11 %%
```

yylval is of type YYSTYPE—**int** by default, but you can change YYSTYPE explicitly by redefining it or implicitly by using the **%union** mechanism described earlier. If you do change the type, be careful to also change any matching **extern** statement in the LEX input file as well.

Note that this mechanism is risky if the lexical analyzer is depending on the parser to have taken some action before a symbol is read. The problem is that the lookahead symbol can be read well in advance of the time that it's shifted—several reductions can occur after reading the symbol but before shifting it. The problem is best illustrated with an example. Say that a grammar for C has two tokens, a **NAME** token that's used for identifiers, and a **TYPE** token that is used for types. A **typedef** statement can cause a string that was formerly treated as a **NAME** to be treated as a **TYPE**. That is, the **typedef** effectively adds a new keyword to the language, because subsequent references to the associated identifier must be treated as if they were **TYPE** tokens.

It is tempting to try to use the symbol table to resolve the problem. When a **typedef** is encountered, the parser creates a symbol-table entry for the new type. The lexical analyzer could then use the symbol table to distinguish **NAMEs** from **TYPEs**. The difficulty here lies in input such as the following:

```

typedef int itype;
itype x;
```

In a bottom-up parser such as the current one, the symbol-table entry for the **typedef** will, typically, not be created until the entire declaration is encountered—and the parser can't know that it's at the end of the declaration until it finds the trailing semicolon. So,

the semicolon is read and shifted on the stack, and the next lookahead (the second `itype`) is read, before doing the reduction that puts the first `itype` into the symbol table—the reduction that adds the table entry happens after the second `itype` is read because the lookahead character must be read to decide to do the reduction. When the scanner looks up the second `itype`, it won't be there yet, so it will assume incorrectly that the second `itype` is a **NAME** token. The moral is that attributes are best put onto the stack during a reduction, rather than trying to put them onto the stack from the lexical analyzer.

E.11.11 Sample Occs Input File

For convenience, the `occs` input file for the expression compiler we've been discussing is shown in Listing E.17.

Listing E.17. `expr.y`—An Expression Compiler (Occs Version)

```

1  %term ID          /* a string of lower-case characters */
2  %term NUM         /* a number */
3
4  %left PLUS        /* + */
5  %left STAR        /* * */
6  %left LP RP        /* ( ) */
7
8  %}
9  #include <stdio.h>
10 #include <ctype.h>
11 #include <malloc.h>
12
13 extern char      *yytext;      /* In yylex(), holds lexeme */
14 extern char      *new_name(); /* declared at bottom of this file */
15
16 typedef char     *stype;       /* Value stack */
17 #define YYSTYPE    stype
18
19 #define YYMAXDEPTH 64
20 #define YYMAXERR   10
21 #define YYVERBOSE
22 %}
23
24 %%
25 /* A small expression grammar that recognizes numbers, names, addition (+),
26 * multiplication (*), and parentheses. Expressions associate left to right
27 * unless parentheses force it to go otherwise. * is higher precedence than +.
28 * Note that an underscore is appended to identifiers so that they won't be
29 * confused with rvalues.
30 */
31
32 s      : e ;
33
34 e      : e PLUS e  { yycode("%s += %s\n", $1, $3); free_name( $3 ); }
35 | e STAR e  { yycode("%s *= %s\n", $1, $3); free_name( $3 ); }
36 | LP e RP   { $$ = $2; }
37 | NUM       { yycode("%s = %s\n", $$ = new_name(), yytext ); }
38 | ID        { yycode("%s = _%s\n", $$ = new_name(), yytext ); }
39 ;
40 %%
41

```



Listing E.17. continued...

```

42  /*-----*/
43  char    *yypstk( ptr )
44  char    **ptr;
45  {
46      /* Yystk is used by the debugging routines. It is passed a pointer to a
47      * value-stack item and should return a string representing that item. In
48      * this case, all it has to do is dereference one level of indirection.
49      */
50
51      return *ptr ? *ptr : "<empty>" ;
52  }
53
54  /*-----*/
55
56  char  *Names[] = { "t0", "t1", "t2", "t3", "t4", "t5", "t6", "t7" };
57  char  **Namep = Names;
58
59  char  *new_name()
60  {
61      /* Return a temporary-variable name by popping one off the name stack. */
62
63      if( Namep >= &Names[ sizeof(Names)/sizeof(*Names) ] )
64      {
65          yyerror("Expression too complex\n");
66          exit( 1 );
67      }
68
69      return( *Namep++ );
70  }
71
72  free_name(s)
73  char  *s;
74  {
75      /* Free up a previously allocated name */
76      *--Namep = s;
77  }
78  /*-----*/
79
80  yy_init_occ()
81  {
82      /* Generate declarations for the rvalues */
83
84      yycode("public word t0, t1, t2, t3;\n");
85      yycode("public word t4, t5, t6, t7;\n");
86  }
87
88  main( argc, argv )
89  char  **argv;
90  {
91      /* Open the input file, using yy_get_args() if we're debugging or
92      * ii_newfile() if not.
93      */
94
95  #ifdef YYDEBUG
96      yy_get_args( argc, argv );
97  #else
98      if( argc < 2 )
99      {
100          fprintf( stderr, "Need file name\n");
101          exit(1);

```



Listing E.17. continued...

```

102      }
103      else if( ii_newfile(argv[1]) < 0 )
104      {
105          fprintf( stderr, "Can't open %s\n", argv[1] );
106          exit(2);
107      }
108  #endif
109      yyparse();
110      exit( 0 );
111  }
```

Listing E.18. *expr.lex*—*LE*X Input File for Expression Compiler

```

1  @@
2  #include "yyout.h"
3  @@
4  @@
5  "+"      return PLUS;
6  "*"      return STAR;
7  "("      return LP;
8  ")"      return RP;
9  [0-9]+   return NUM;
10 [a-z]+   return ID;
11 %%
```

E.11.12 Hints and Warnings

- Though input defaults to standard input in production mode (as compared to debug mode), the input routines really expect to be working with a file. The default, standard-input mode is intended for use in pipes, not for interactive input with a human being. This expectation can produce unexpected consequences in an interactive situation. Since the input is always one token ahead of the parser, the parser's actions can appear to be delayed by one token. This delay is most noticeable at end of file, because the last token in the input stream isn't processed until an explicit end-of-file marker is read. If you're typing the input at the keyboard, you'll have to supply this marker yourself. (Use *Ctrl-D* with UNIX; the two-character sequence *Ctrl-Z Enter* under MS-DOS).
- Though you can redefine YYPRIVATE to make various global **static** variables public for the purpose of debugging, you should never modify any of these global variables directly.
- Once the grammar is working properly, make changes very carefully. Very small changes in a grammar can introduce masses of shift/reduce and reduce/reduce conflicts. You should always change only one production at a time and then remake the tables. Always back up the current input file before modifying it so that you can return to a working grammar if you mess up (or use a version-control system like SCCS).
- Avoid ϵ productions and imbedded actions—they tend to introduce shift/reduce conflicts into a grammar. If you must introduce an imbedded production, try to put it immediately to the right of a terminal symbol.
- Avoid global variables in the code-generation actions. The attribute mechanism should be used to pass all information between productions if at all possible

(sometimes it's not). Grammars are almost always recursive. Consequently, you'll find that global variables tend to be modified at unexpected times, often destroying information that you need for some subsequent action. Avoiding global variables can seem difficult. You have to work harder to figure out how to do things—it's like writing a program that uses subroutine return values only, without global variables or subroutine arguments. Nonetheless, it's worth the effort in terms of reduced development time. All of the global-variable-use issues that apply to recursive subroutines apply to the action code in a production.

- The assignment in the default action (`$$=$1`) is made before the action code is executed. Modifying `$1` inside one of your own actions will have no effect on the value of `$$`. You should modify `$$` itself.
- If you are using the `-a` and `-p` switches to split the parser into two files, remember that actions imbedded in a production actually modify the grammar. If you add or move such an action, you must remake the tables. You can add or remove actions at the far right of a production without affecting the tables, however.

Occs is not **yacc**, and as a consequence, many hacks found in various books that discuss the UNIX utilities must be avoided when using **occs**:

- Because of the way that the input is processed, it's not safe to modify the lexeme from the parser or to do any direct input from the parser. All tokens should be returned from L^EX in an orderly fashion. You must use `yytext`, `yylineno`, and `yylen` to examine the lexeme. It's risky for code in the parser to modify these variables or to call any of the `ii_` input routines used by L^EX. The problem here is that **occs** and **LLama** both read one token ahead—a second, lookahead token will already have been read before any action code is processed. The token in `yytext` is the current token, not the lookahead token.
- By the same token (so to speak) you should never modify the **occs** or **LLama** value stack directly, always use the dollar-attribute mechanism—`$$`, `$1`, and so on—to do so. The contents of the `yyval`, which has the same type as a stack element, is shifted onto the value stack when a state representing a token is shifted onto the state stack, and you can use this variable to shift an attribute for a token. (Just assign a value to `yyval` before returning the token). It is not possible for code in a L^EX-generated lexical analyzer to modify the value stack directly, as is done in some published examples of how to use the UNIX utilities. Use `yyval`.
- The **occs** error-recovery mechanism is completely automatic. Neither the **yacc** error token, nor the `yyerrok` action is supported by **occs**. The `error` token can be removed from all **yacc** grammars. Similarly, all `yyerrok` actions can be deleted. If a **yacc** production contains nothing but an error token and optional action on its right-hand side, the entire production should be removed (don't just delete the right-hand side, because you'll introduce a hitherto nonexistent ϵ production into the grammar).
- The **occs** start production may have only one right-hand side. If a **yacc** grammar starts like this:

```
baggins : frodo
          | bilbo
          ;
```

add an extra production at the very top of the **occs** grammar (just after the `%%`):

```
start : baggins ;
```

E.12 LLama

The remainder of this appendix describes the **LLama**-specific parts of the compiler compiler. The main restriction in using **LLama** is that the input grammar must be LL(1). **LLama** grammars are, as a consequence, harder to write than **occ**s grammars. On the other hand, a **LLama**-generated parser will be both smaller and faster than an **occ**s parser for an equivalent grammar.

E.12.1 Percent Directives and Error Recovery

The % directives supported by **LLama** are summarized in Table E.8. **LLama** supports one % directive over and above the standard directives. The %synch is placed in the definitions section of the input file—use it to specify a set of *synchronization tokens* for error recovery. A syntax error is detected when the top-of-stack symbol is a terminal symbol, and that symbol is not also the current lookahead symbol. The error-recovery code does two things: it pops items off the stack until the top-of-stack symbol is a token in the synchronization set, and it reads tokens from input until it finds the same token that it just found on the stack, at which point it has recovered from the error. If the parser can't find the desired token, or if no token in the synchronization set is also on the stack, then the error is unrecoverable and the parse is terminated with an error flag. `yyparse()` usually returns 1 in this case, but this action can be changed by redefining `YYABORT`. (See table E.2 on page 840).

The %synch directive, synchronization tokens.

Changing the action taken by the parser on an error.

Table E.8. LLama % Directives and Comments

Directive	Description
%%	Delimits the three sections of the input file.
% {	Starts a code block. All lines that follow, up to a %} are written to the output file unchanged.
% }	Ends a code block.
%token	Defines a token.
%term	A synonym for %token.
/* */	C-like comments are recognized—and ignored—by occ s, even if they're outside of a %{ %} delimited code block.
%synch	Define set of synchronization tokens.

Several tokens can be listed in the %synch directive. Good choices for synchronization symbols are tokens that end something. In C, for example, semicolons, close parentheses, and close braces are reasonable selections. You'd do this with:

```
%term SEMI CLOSE_PAREN CLOSE_CURLY
. .
%synch SEMI CLOSE_PAREN CLOSE_CURLY
```

E.12.2 Top-Down Attributes

The **LLama** value stack and the \$ attribute mechanism are considerably different from the one used by **occ**s. **LLama** uses the top-down attribute-processing described in Chapter Four. Attributes are referenced from within an action using the notation \$\$, \$1, \$2, and so forth. \$\$ is used in an action to reference the attribute that was attached to the nonterminal on the left-hand side before it was replaced. The numbers can be used to reference attributes attached to symbols to the right of the current action in the grammar. The number indicates the distance from the action to the desired symbol. (\$0 references

LLama attributes, \$\$, \$1, etc.

the current action's own attributes.) For example, in

```
stmt :           {$1=$2=new_name();} expr {free_name($0);} SEMI stmt ;
```

the \$1 in the left action modifies the attribute attached to expr, the \$2 references the attribute attached to the second action, which uses \$0 to get that attribute. \$\$ references the attribute attached to the left-hand side in the normal way. Attributes flow across the grammar from left to right.

E.12.3 The LLama Value Stack

Typing LLama's value stack.

The **LLama** value stack is a stack of structures, as was described in Chapter Four. The structure has two fields, defined in the **LLama**-generated parser as follows:

```
typedef struct      /* Typedef for value-stack elements.      */
{
    YYSTYPE left; /* Holds value of left-hand side attribute. */
    YYSTYPE right; /* Holds value of current-symbol's attribute.*/
}
yyvstype;
```

The YYSTYPE structure is defined as an **int** by default, but you can redefine it in the definitions section of the input file as follows:

```
%{
typedef char *stack_type; /* Use stack of character pointers. */
#define YYSTYPE stack_type
%}
%%
```

In this case, the dollar attribute will be of the same type as YYSTYPE. That is \$\$, \$1, and so forth, reference the character pointer. You can use *\$1 or \$1[2] to access individual characters in the string. If the stack is a stack of structures or unions, as in:

```
%{
typedef union
{
    int      integer;
    char *string;
}
stack_type; /* Use stack of character pointers */
#define YYSTYPE stack_type
%}
%%
```

You can access individual fields like this: \$\$.integer or \$1.string.

Initializing the value stack with yy_init_llama().

The initialization subroutine for the **LLama**-generated parser, **yy_init_llama()**, is called after the stack is initialized but before the first lexeme is input—the start symbol will have been pushed onto the parse stack, and garbage onto the value stack. The initialization routine is passed a pointer to the garbage entry for the pushed start symbol. The default routine in *l.lib* does nothing, but you can use your own **yy_init_llama(p)** to provide an attribute for the goal symbol so that subsequent replacements won't inherit garbage. Your replacement should be put at the bottom of the **LLama** input file and should look something like this:

```
yy_init_llama( p )
yyvstype *p;
{
    p->left = p->right = initial attribute value for goal symbol;
}
```

The yyvstype type is a two-part structure used for the value-stack items described earlier.

The yypstk() routine used to print value-stack items in the debugging environment is passed a pointer to a value-stack item of type (a pointer to a yyvstype structure) and a string representing the symbol. You could print character-pointer attributes like this:

```
%{
typedef char *stack_type; /* Use stack of character pointers.*/
#define YYSTYPE stack_type
%
%
...
%%
char *yypstk(tovs, tods) /* Print attribute stack contents.*/
yyvstype *tovs;
char *symbol;
{
    static char buf[64];
    sprintf(buf, "[%0.30s,%0.30s]", tovs->left, tovs->right);
    return buf;
}
```

Printing the value-stack,
yypstk(),yyvstype.

E.12.4 The *llout.sym* File

LLama produces a symbol table file if the *-s* or *-S* switch is specified on the command line. Listing E.19 shows the symbol table produced by the input file at the end of this appendix.

Generating the symbol-table file, -s -S.

The first part of the file is the nonterminal symbols. Taking expr as characteristic, the entry looks like this:

```
expr (257)
FIRST : NUM_OR_ID LP
FOLLOW: RP SEMI
2: expr -> term expr' .....SELECT: NUM_OR_ID LP
```

The 257 is the symbol used to represent an expr on the parse stack, the next two lines are the FIRST and FOLLOW sets for expr. (These are present only if *-S* was used to generate the table.) The next lines are all productions that have an expr on their left-hand side. The number (2, here) is the production number, and the list to the right is the LL(1) selection set for this production. The production number is useful for setting breakpoint on application of that production in the debugging environment.

If the production contains an action, a marker of the form {*N*} is put into the production in place of the action. All these action symbols are defined at the bottom of *llout.sym* (on lines 51 to 57 of the current file). Taking {0} as characteristic:

```
{0} 512, line 42 : {$1=$2=newname();}
```

The 512 is the number used to represent the action on the parse stack, the action was found on line 42 of the input file, and the remainder of the line is the first few characters of the action itself.

The middle part of the symbol table just defines tokens. The numbers here are the same numbers that are in *llout.h*, and these same values will be used to represent a token on the parse stack.

E.12.5 Sample LLama Input File

Listing E.20 is a small expression compiler in LLama format.

Listing E.19. A LLama Symbol Table

```

1 ----- Symbol table -----
2
3 NONTERMINAL SYMBOLS:
4
5 expr (257)
6   FIRST : NUM_OR_ID LP
7   FOLLOW: RP SEMI
8     2: expr -> term expr' .....SELECT: NUM_OR_ID LP
9
10 expr' (259)
11   FIRST : PLUS <epsilon>
12   FOLLOW: RP SEMI
13     3: expr' -> PLUS {2} term {3} expr' ....SELECT: PLUS
14     4: expr' -> .....SELECT: RP SEMI
15
16 factor (260)
17   FIRST : NUM_OR_ID LP
18   FOLLOW: PLUS TIMES RP SEMI
19     8: factor -> NUM_OR_ID {6} .....SELECT: NUM_OR_ID
20     9: factor -> LP expr RP .....SELECT: LP
21
22 stmt (256) (goal symbol)
23   FIRST : NUM_OR_ID LP <epsilon>
24   FOLLOW: $
25     0: stmt -> .....SELECT: $
26     1: stmt -> {0} expr {1} SEMI stmt .....SELECT: NUM_OR_ID LP
27
28 term (258)
29   FIRST : NUM_OR_ID LP
30   FOLLOW: PLUS RP SEMI
31     5: term -> factor term' .....SELECT: NUM_OR_ID LP
32
33 term' (261)
34   FIRST : TIMES <epsilon>
35   FOLLOW: PLUS RP SEMI
36     6: term' -> TIMES {4} factor {5} term' ..SELECT: TIMES
37     7: term' -> .....SELECT: PLUS RP SEMI
38
39 TERMINAL SYMBOLS:
40
41 name           value
42 LP              4
43 NUM_OR_ID      3
44 PLUS            1
45 RP              5
46 SEMI            6
47 TIMES           2
48
49 ACTION SYMBOLS:
50
51 {0} 512, line 42 : {$1=$2=newname();}
52 {1} 513, line 42 : {freename($0);}
53 {2} 514, line 48 : {$1=$2=newname();}
54 {3} 515, line 49 : { yycode("%s+=%s\\n", $$, $0); freename($0); }
55 {4} 516, line 56 : {$1=$2=newname();}
56 {5} 517, line 57 : { yycode("%s*=%s\\n", $$, $0); freename($0); }
57 {6} 518, line 61 : { yycode("%s=%0.*s\\n", $$, yytext); }

```

Listing E.20. *expr.lma*— An Expression Compiler (Llama Version)

```

1 %term PLUS          /* + */          */
2      TIMES         /* * */          */
3 %term NUM_OR_ID    /* a number or identifier */ */
4 %term LP            /* ( */          */
5 %term RP            /* ) */          */
6 %term SEMI          /* ; */          */
7
8 %}
9  -----
10 /* Rvalue names are stored on a stack. name() pops a name off the stack and
11  * freename(x) puts it back. A real compiler would do some checking for
12  * stack overflow here but there's no point in cluttering the code for now.
13 */
14
15 char *Namepool[] =
16 {
17     "t0", "t1", "t2", "t3", "t4", "t5", "t6", "t7", "t8", "t9"
18 };
19 char **Namep = Namepool ;
20
21 char *newname()           { return(*Namep++); }
22 char *freename( char *x) { return(*--Namep = x); }
23
24 extern char *yytext;
25 extern int yyleng;
26
27 #define YYSTYPE char*
28 %}
29
30 %synch SEMI RP
31
32 %%
33 /* A small expression grammar that recognizes numbers, names, addition (+),
34  * multiplication (*), and parentheses. Expressions associate left to right
35  * unless parentheses force it to go otherwise. * is higher precedence than +
36 */
37
38 stmt   : /* eps */
39   | {$1=$2=newname();} expr {freename($0);} SEMI stmt
40   ;
41
42 expr   : term expr'
43   ;
44
45 expr'  : PLUS {$1=$2=newname();} term
46           { yycode("%s+=%s\n", $$, $0); freename($0); } expr'
47   | /* epsilon */
48   ;
49
50 term   : factor term'
51   ;
52
53 term'  : TIMES {$1=$2=newname();} factor
54           { yycode("%s*=%s\n", $$, $0); freename($0); } term'
55   | /* epsilon */
56   ;
57

```



Listing E.20. continued...

```
58 factor  :      NUM_OR_ID { yycode("%s=%0.*s\n", $$, yyleng, yytext); }
59      |      LP expr RP
60      ;
61 %%
62 /*-----*/
63
64 yy_init_llama( p )
65 yyvstype *p;
66 {
67     p->left = p->right = "-";
68 }
69
70 char *yyppstk(tovs, tod)
71 yyvstype *tovs;
72 char      *tods;
73 {
74     static char buf[128];
75     sprintf(buf, "[%s,%s]", tovs->left, tovs->right);
76     return buf;
77 }
78
79 main( argc, argv )
80 char      **argv;
81 {
82     yy_get_args( argc, argv );
83     yyparse();
84     exit( 0 );
85 }
```

A C-code Summary

This appendix summarizes all the C-code directives described in depth in Chapter Six. Many of the tables in that chapter are also found here.

Figure F.1. The C-code Virtual Machine

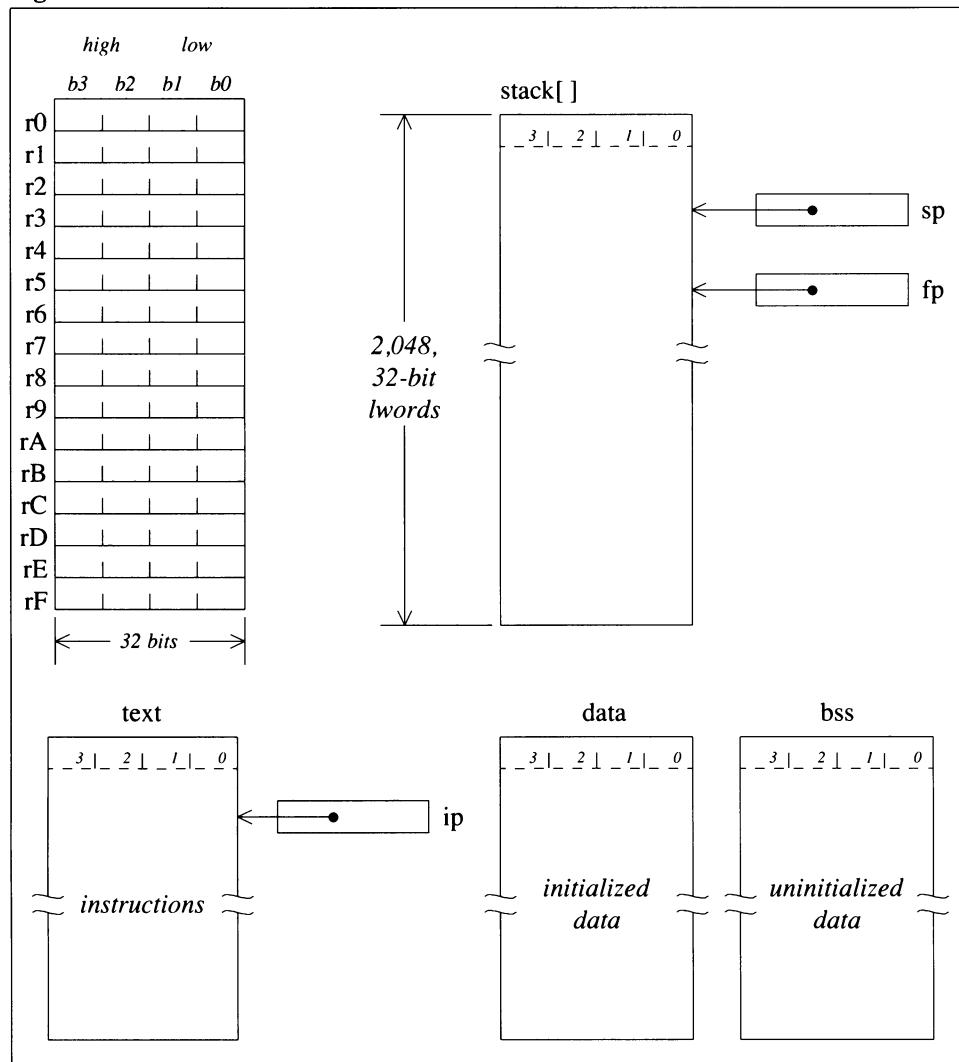


Table F.1. Registers

<code>rN</code>	General-purpose register (<code>r0, r1, ..., r9, rA, ..., rF</code>).
<code>sp</code>	Stack pointer.
<code>fp</code>	Frame pointer.
<code>ip</code>	Instruction pointer.
<code>rN.pp</code>	Access 32-bit register as pointer to something. Must use an addressing mode to specify type of object pointed to.
<code>rN.w.high</code>	Access word in high 16-bits of register.
<code>rN.w.low</code>	Access word in low 16-bits of register.
<code>rN.b.b3</code>	Access most significant byte of register (MSB).
<code>rN.b.b2</code>	Access low byte of high word.
<code>rN.b.b1</code>	Access high byte of low word.
<code>rN.b.b0</code>	Access least significant byte of register (LSB).

Table F.2. Types, Storage Classes, and Declarations

Types	
<code>byte</code>	8-bit
<code>array</code>	(alias for <code>byte</code> —used to declare pointers to arrays)
<code>record</code>	(alias for <code>byte</code> —used to declare structures)
<code>word</code>	16-bit
<code>lword</code>	32-bit
<code>ptr</code>	generic pointer
Storage classes	
<code>private</code>	Space is allocated for the variable, but the variable can not be accessed from outside the current file.
<code>public</code>	Space is allocated for the variable, and the variable can be accessed from any file in the current program. It is illegal for two <code>public</code> variables to have the same name, even if they are declared in separate files.
<code>external</code>	Space for this variable is allocated elsewhere. There must be a <code>public</code> or <code>common</code> definition for the variable in some other file.
<code>common</code>	Space for this variable is allocated by the linker. If a variable with a given name is declared <code>common</code> in one module and <code>public</code> in another, then the <code>public</code> definition takes precedence. If there are nothing but <code>common</code> definitions for a variable, then the linker allocates space for that variable in the <code>bss</code> segment.
Declarations	
<code>class type name;</code>	Variable of indicated <code>type</code> and storage <code>class</code> .
<code>class type name [size];</code>	Array of indicated <code>type</code> and storage <code>class</code> ; <code>size</code> is optional if initialized with C-style initializer.
All declarations in the <code>data</code> segment <u>must</u> be initialized using a C-style initializer. Nonetheless, pointers may not be initialized with string constants, only with the address of a previously-declared variable. Declarations in the <code>bss</code> segment <u>may not</u> be initialized.	

Table F.3. Converting C Storage Classes to C-code Storage Classes

	static	not static	
		definition	declaration (extern or prototype)
Subroutine (in <i>text</i> segment).	private	public	external
Uninitialized variable (in <i>bss</i> segment).	private	common	
Initialized variable (in <i>data</i> segment).	private	public	

Table F.4. Addressing Modes

Mode	Example	Notes		
immediate	10	Decimal number. Use leading 0x for hex, 0 for octal.		
direct	x, r0.1	Contents of variable or register.		
indirect	B(p)	byte whose address is in <i>p</i> .		
	W(p)	word whose address is in <i>p</i> .		
	L(p)	lword whose address is in <i>p</i> .		
	P(p)	ptr whose address is in <i>p</i> .		
	BP(p)	byte pointer whose address is in <i>p</i> .		
	WP(p)	word pointer whose address is in <i>p</i> .		
	LP(p)	lword pointer whose address is in <i>p</i> .		
	PP(p)	ptr pointer whose address is in <i>p</i> .		
double indirect	*BP(p)	byte pointed to by pointer whose address is in <i>p</i> .		
	*WP(p)	word pointed to by pointer whose address is in <i>p</i> .		
	*LP(p)	lword pointed to by pointer whose address is in <i>p</i> .		
	*PP(p)	word pointed to by pointer whose address is in <i>p</i> .		
based indirect	P(p+N)	ptr at byte offset <i>N</i> from address in <i>p</i> .		
	B(p+N)	word at byte offset <i>N</i> from address in <i>p</i> .		
	W(p+N)	lword at byte offset <i>N</i> from address in <i>p</i> .		
	L(p+N)	lword at byte offset <i>N</i> from address in <i>p</i> .		
	LP(p+n)	lword pointer at byte offset <i>N</i> from address in <i>p</i> .		
	*LP(p+n)	lword pointed by pointer at byte offset <i>N</i> from address in <i>p</i> .		
... . . .				
effective address	&name	Address of variable or first element of array.		
	&W(p+N)	Address of word at offset + <i>n</i> from the pointer <i>p</i> .		
	The effective-address modes can also be used with other indirect modes. (See Table F.5).			
A generic pointer, <i>p</i> , is a variable declared ptr or a pointer register (rN, pp, fp, sp).				
<i>N</i> is any integer: a number, a numeric register (r0.w.low), or a reference to a byte, word, or lword variable.				

Table F.5. Combined Indirect and Effective-Address Modes

Syntax:	Evaluates to:
$\&p$ $\&W(\&p)$ $\&WP(\&p)$	address of the pointer
p $\&W(p)$ $WP(\&p)$	contents of the pointer itself
$W(p)$ $*WP(\&p)$	contents of the word whose address is in the pointer

Table F.6. Arithmetic Operators

$d = s$	assignment
$d += s$	addition
$d -= s$	subtraction
$d *= s$	multiplication
$d /= s$	division
$d \%= s$	modulus division
$d ^= s$	bitwise XOR
$d &= s$	bitwise AND
$d = s$	bitwise OR
$d <<= s$	Shift d to left by s bits
$d >>= s$	Shift d to right by s bits (arithmetic)
$d == s$	d = twos complement of s
$d =~ s$	d = ones complement of s
$d =& s$	d = effective address of s
$lrs(d,n)$	logical right shift of d by n bits. Zero fill in left bits.

Table F.7. Test Directives

Directive:	Execute following line if:	
EQ(a, b)	$a = b$	
NE(a, b)	$a \neq b$	
LT(a, b)	$a < b$	
LE(a, b)	$a \leq b$	
GT(a, b)	$a > b$	
GE(a, b)	$a \geq b$	
U_LT(a, b)	$a < b$	(unsigned comparison)
U_LE(a, b)	$a \leq b$	(unsigned comparison)
U_GT(a, b)	$a > b$	(unsigned comparison)
U_GE(a, b)	$a \geq b$	(unsigned comparison)
BIT(b, s)	bit b of s is set to 1 (bit 0 is the low bit).	

Table F.8. Other Directives

<code>SEG(text)</code>	Change to code segment.
<code>SEG(data)</code>	Change to initialized-data segment.
<code>SEG(bss)</code>	Change to uninitialized-data segment.
<code>PROC (name)</code>	Begin code for subroutine <i>name</i> .
<code>ENDP (name)</code>	End code for subroutine <i>name</i> .
<code>ALIGN (type)</code>	Next declared variable is aligned as if it were the indicated type.
<code>ext_low(reg);</code>	Duplicate high bit of <code>reg.b.b0</code> in all bits of <code>reg.b.b1</code> .
<code>ext_high(reg);</code>	Duplicate high bit of <code>reg.b.b2</code> in all bits of <code>reg.b.b3</code> .
<code>ext_word(reg);</code>	Duplicate high bit of <code>reg.w.low</code> in all bits of <code>reg.w.high</code> .
<code>push(x);</code>	Push <i>x</i> onto the stack.
<code>x = pop(type);</code>	Pop an object of the indicated type into <i>x</i> .
<code>call(ref);</code>	Call subroutine. <i>ref</i> can be a subroutine name or a pointer.
<code>ret();</code>	Return from current subroutine.
<code>link(N);</code>	Set up stack frame: <code>push(fp), fp=sp; sp-=N</code> .
<code>unlink();</code>	Discard stack frame: <code>sp=fp, fp=pop(ptr)</code> .
<code>pm()</code>	(subroutine) Print all registers and top few stack items.
<code>_main</code>	is mapped to <code>main</code> .

Directives whose names are in all upper case must not be followed by a semicolon when used. All other directives must be found inside a subroutine delimited with PROC () and ENDP () directives. SEG () may not be used in a subroutine.

Table F.9. Supported Preprocessor Directives

<code>#line line-number "file"</code>	Generate error messages as if we were on line <i>line-number</i> of the indicated <i>file</i> .
<code>#define NAME text</code>	Define macro.
<code>#define NAME (x) text</code>	Define macro with arguments.
<code>#undef NAME</code>	Delete macro definition.
<code>#ifdef NAME</code>	If macro is defined, compile everything up to next <code>#else</code> or <code>#endif</code> .
<code>#if constant-expression</code>	If <i>constant-expression</i> is true, compile everything up to next <code>#else</code> or <code>#endif</code> .
<code>#endif</code>	End conditional compilation.
<code>#else</code>	If previous <code>#ifdef</code> or <code>#if</code> is false, compile everything up to the next <code>#endif</code> .
<code>#include <file></code>	Replace line with contents of file (search system directory first, then current directory).
<code>#include "file"</code>	Replace line with contents of file (search current directory).

Table F.10. Return Values: Register Usage

Type	Returned in:
<code>char</code>	<code>rF.w.low</code> (A <code>char</code> is always promoted to <code>int</code> .)
<code>int</code>	<code>rF.w.low</code>
<code>long</code>	<code>rF.l</code>
<code>pointer</code>	<code>rF.pp</code>

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WP	6.9(49) <i>virtual.h</i> , 465	yyhook_a()	4.13(1) <i>yyhook_a.c</i> , 264
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 YY_PUSH() 5.12(177) *yyout.c*, 389
 YY_PUSH() E.20(27) *expr.lma*, 887
 YY_PUSH() E.8(17) *expr.y*, 857
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Errata: *Compiler Design in C*

This document is a list of typos and corrections that need to be made to *Compiler Design in C*, Allen Holub, 1990 (as of September 11, 1997). The corrections marked “Disk only” represent changes made to the files on the distribution disk that either don’t affect the code in any significant way or are too big to insert into the book. All these will eventually be incorporated into the second edition, if there is one, but they won’t be put into subsequent printings of the first edition. There are also many trivial changes that have been made to the code on the disk to make it more robust: Function prototypes have been added here and there, as have includes for .h files that contain prototypes, all .h files have been bracketed with statements like:

```
#ifndef __FILE_EXT /* file name with _ for dot */
#define __FILE_EXT
...
#endif
```

to make multiple inclusions harmless, and so forth. These changes allow the code to compile under Borland C++ as well as Microsoft C and BSD UNIX. None of these trivial changes are documented here.

1, 2 (1.02)

The printings in which the error occurs and the software version in which the change was made are identified in the margin at the head of each entry. For example, the numbers next to the current paragraph indicate that the typo is found in both the first and second printings, and that the bug was fixed in software version 1.02. Determine your printing by looking at the back of the flyleaf, where you’ll find a list of numbers that looks something like this: 10 9 8 7 6 5 4 3. The smallest number in the list is the printing.

Page xvi - Seventh line from the bottom. Change “No credit cards” to “No purchase orders or credit cards.” The last two lines of the paragraph should read:

must add local sales tax. **No purchase orders or credit cards** (sorry). A Macintosh version will be available eventually. Binary site licenses are available for educational institutions.

Page xvi - Last line. Internet can now be accessed from CompuServe. Add the Compuserve/internet address >INTERNET:holub@violet.berkeley.edu to the parenthesized list at the end of the paragraph. A replacement paragraph follows:

The code in this book is bound to have a few bugs in it, though I’ve done my best to test it as thoroughly as possible. The version distributed on disk will always be the most recent. If you find a bug, please report it to me, either at the above address or electronically. My internet address is *holub@violet.berkeley.edu* CompuServe users can access internet from the email system by prefixing this address with >INTERNET: —type help internet for information. My UUCP address is ...!ucbvax!violet!holub.

Page xvi - 15 lines from the bottom. Change the phone number to:

(510) 540-7954

Page xviii - Line 7, change “that” to “than”. The replacement line follows:

primitives. It is much more useful than **pic** in that you have both a WYSIWYG capability ‘

‘
Page 8 - The line that starts “**JANE verb object**” in the display at the bottom of the page is repeated. Delete the first one. A replacement display follows:

<i>sentence</i>	apply <i>sentence</i> → <i>subject predicate</i> to get:
<i>subject predicate</i>	apply <i>subject</i> → <i>noun</i> to get:
<i>noun predicate</i>	apply <i>noun</i> → JANE to get:
JANE predicate	apply <i>predicate</i> → <i>verb object</i> to get:
JANE verb object	apply <i>verb</i> → SEES to get:
JANE SEES object	apply <i>object</i> → <i>noun op_participle</i> to get:
JANE SEES noun opt_participle	apply <i>noun</i> → SPOT to get:
JANE SEES SPOT opt_participle	apply <i>opt_participle</i> → <i>participle</i> to get:
JANE SEES SPOT participle	apply <i>participle</i> → RUN to get:
JANE SEES SPOT RUN	done—there are no more nonterminals to replace

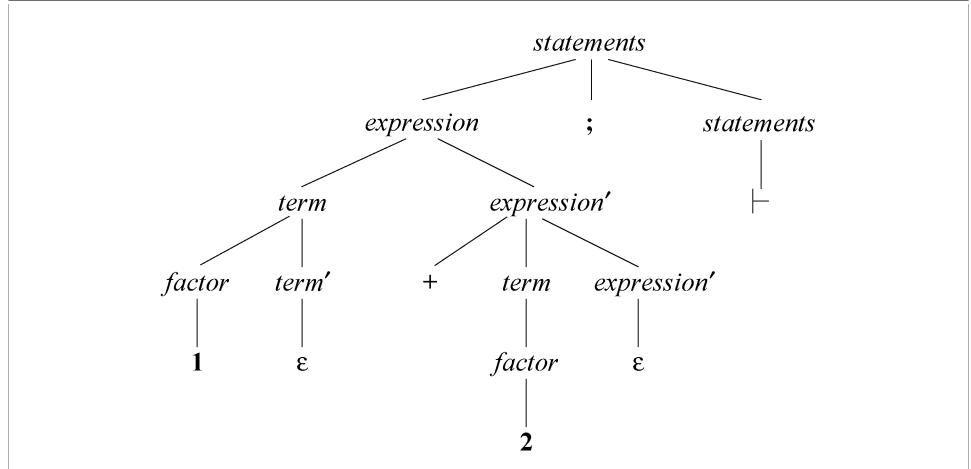
Page 11 - Table 1.1, line 1 should read: *statements*→**EOF**. A replacement table follows:

Table 1.1. A Simple Expression Grammar

1.	<i>statements</i>	→	EOF
2.			<i>expression ; statements</i>
3.	<i>expression</i>	→	<i>expression + term</i>
4.			<i>term</i>
5.	<i>term</i>	→	<i>term * factor</i>
6.			<i>factor</i>
7.	<i>factor</i>	→	number
8.			(<i>expression</i>)

Page 12 - Figure 1.6. Replace the figure with the following one:

Figure 1.6. A Parse of 1+2



Page 16 - Listing 1.2, delete line 46. (Replace it with a blank line.)

Page 18 - Replace the untitled table just under Listing 1.4 with the following (only the first line has been changed):

1.	$statements \rightarrow expression ; eoi$
2.	$ \quad expression ; statement$
3.	$expression \rightarrow term expression'$
4.	$expression' \rightarrow + term expression'$
5.	$ \quad \epsilon$
6.	$term \rightarrow factor term'$
7.	$term' \rightarrow * factor term'$
8.	$ \quad \epsilon$
9.	$factor \rightarrow num_or_id$
10.	$ \quad (expression)$

Page 24 - Line 110 of Listing 1.7 should read $*p++ = tok;$

110

$*p++ = tok;$

Page 26 - Change the caption to Figure 1.11 as follows:

Figure 1.11. A Subroutine Trace of 1+2*3+4 (Improved Parser)

Page 27 - Change Line 42 of Listing 27 to `tempvar2 = term();`

```
42      tempvar2 = term();
```

Page 27 - Replace the word `temporary` in the code part of lines 19, 26, and 42 of Listing 1.9 with the word `tempvar`. These three lines should now read:

```
19      tempvar = expression()
26      freename( tempvar );
```

Page 36 - Replace Lines 16 and 17 of Listing 2.1 with the following (I've added a few parentheses):

```
16 #define get(stream)  (Pbackp < &Pbackbuf[SIZE] ? *Pbackp++ : getc(stream) )
17 #define unget(c)     (Pbackp <= Pbackbuf ? -1 : (*--Pbackp=(c)) )
```

Page 44 - Replace line 186 of Listing 2.5 with the following line:

```
186      Next = sMark = eMark = END - 1;  paMark = NULL;
```

Page 55 - Eleven lines from bottom. The display that says `([^a-z] | \en)` should read as follows:

`([^a-z] | \n)`

Page 41 - Modify line 113 of Listing 2.3 to the following:

```
113      eMark     = END;  pMark = NULL;
```

Page 46 - Replace lines 299 and 300 of Listing 2.6 with the following lines:

```
299     int need, /* Number of bytes required from input. */
300     got; /* Number of bytes actually read. */
```

Page 57 - Line six should read as follows:

causes a transition to State 1; from State 1, an e gets the machine to State 2, and an i ‘

Page 57 - Eleventh line from the bottom should read as follows:

```
next_state = Transition_table[ current_state ][ input_character ];
```

Page 57 - Last line of second paragraph should read as follows (delete the "s"):

r, or e from State 0) are all implied transitions to a special implicit error state. ‘

Page 63 - Listing 2.11, lines 2 and 3: remove the semicolons.

Page 68 - The first line beneath the Figure should read as follows (the Figure and Listing numbers are wrong):

table is shown in Figure 2.6 and in Listing 2.15. The YY_cmap [] array is indexed by ‘

Page 72 - The first display (second and third lines) should read as follows:

```
[‘0’,2] [‘1’,2] [‘2’,2] [‘3’,2] [‘4’,2]
[‘5’,2] [‘6’,2] [‘7’,2] [‘8’,2] [‘9’,2] [‘e’,5]
```

Page 73 - The first line of the third paragraph calls out the wrong line number. It should read as follows:

The YYERROR() macro on line 151 of Listing 2.19 prints internal error messages. ‘

Page 73 - First paragraph, lines two and four. Both references to Listing 2.18 should be to Listing 2.19. A replacement first paragraph follows:

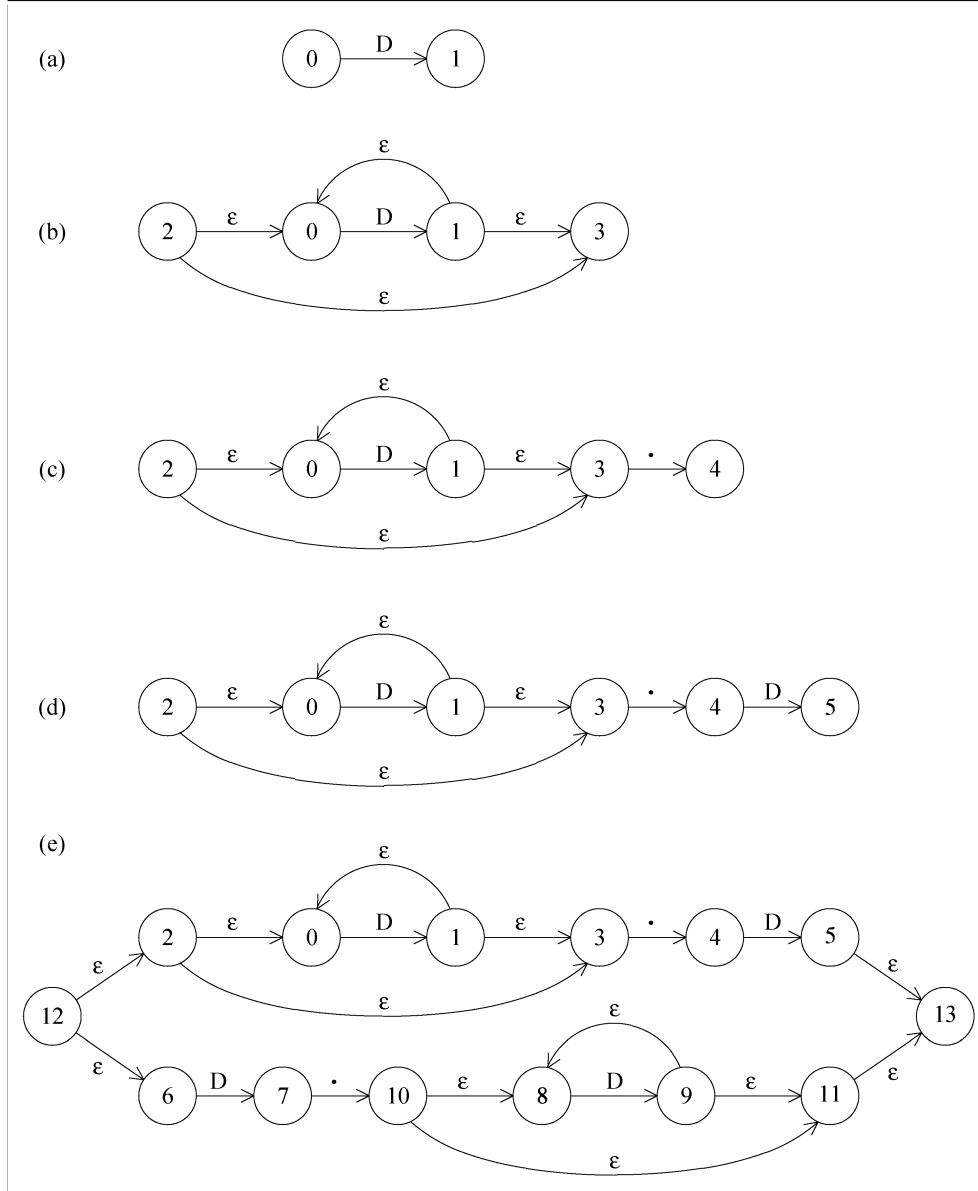
The remainder of lexyy.c file is the actual state-machine driver, shown in Listing 2.19. The first and last part of this listing are the second and third parts of the Ctrl-L-delimited template file discussed earlier. The case statements in the middle (on lines 287 to 295 of Listing 2.19) correspond to the original code attached to the regular expressions in the input file and are generated by LeX itself.

Page 76 - Listing 2.19, line 166. Some code is missing. Replace line 166 with the following line:

Page 78 - Listing 2.19, lines 292 and 294. Align the F in FCON under the I in ICON (on line 288).

Page 84 - Figure 2.13(e). The arrows should point from states 5 and 11 to state 13. Here is a replacement figure:

Figure 2.13. Constructing an NFA for $(D^* \setminus .D | D \setminus .D^*)$



Page 85 - Third line from the bottom (which starts “For example”) should read:

For example, in a machine with a 16-bit int, the first two bytes of the string are the ‘‘

Page 86 - Listing 2.21. Change the number 512 on line 28 to 768.

768

Page 91 - Change the definition of STACK_USED on line 93 of Listing 2.25 to ((int)(Sp-Sstack)+1). A replacement line follows:

```
93 #define STACK_USED() ((int)(Sp-Sstack) + 1) /* slots used */
```

Page 92 - Change line 167 of Listing 2.25 to the following:

```
167 if( textp >= (char *)strings + (STR_MAX-1) )
```

Page 95 - Listing 2.26. Replace lines 280 to 283 with the following lines:

```
280 *p = '\0'; /* Overwrite close brace. { */
281 if( !(mac = (MACRO *) findsym( Macros, *namep )) )
282     parse_err( E_NOMAC );
283 *p++ = '}'; /* Put the brace back. */
```

Page 104 - Second paragraph, first four lines (above the first picture) should read:

Subroutines expr() and cat_expr() are in Listing 2.30. These routines handle the binary operations: | (OR) and concatenation. I’ll show how expr() works by watching it process the expression A|B. The cat_expr() call on line 621 creates a machine that recognizes the A:

Page 121 - Listing 2.38. Replace lines 218 to 255 with the following:

```

start_dfastate = newset();           /* 2 */
ADD( start_dfastate, sstate );
if( !e_closure( start_dfastate, &accept, &anchor ) )
{
fprintf(stderr, "Internal error: State machine is empty\n");
exit(1);
}
current = newset();           /* 3 */
ASSIGN( current, start_dfastate );

/* Now interpret the NFA: The next state is the set of all NFA states that
 * can be reached after we've made a transition on the current input
 * character from any of the NFA states in the current state. The current
 * input line is printed every time an accept state is encountered.
 * The machine is reset to the initial state when a failure transition is
 * encountered.
*/
while( c = nextchar() )
{
next = e_closure( move(current, c), &accept, &anchor );
if( accept )
{
    printbuf();           /* accept */
    if( next );
        delset( next );           /* reset */
        ASSIGN( current, start_dfastate );
}
else
{
    /* keep looking */
    delset( current );
    current = next;
}
}
delset( current );           /* Not required for main, but you'll */
delset( start_dfastate );           /* need it when adapting main() to a */
                                    /* subroutine. */
}
#endif

```

Page 122 - First display, which starts with ‘ε-closure({12})’, should read as follows:

$$\epsilon\text{-closure}(\{0\}) = \{0, 1, 3, 4, 5, 12\} \quad (\text{new DFA State } 0)$$

Page 123 - Display on lines 6–9. Second line of display, which now reads ‘ε-closure({7,11}) = {9, 11, 13, 14},’ is wrong. The entire display should read as follows:

DFA State 7 = {11}	
$\epsilon\text{-closure}(\{11\})$	= {9, 11, 13, 14}
$\text{move}(\{9, 11, 13, 14\}, .)$	= \emptyset
$\text{move}(\{9, 11, 13, 14\}, \mathbf{D})$	= {11} (existing DFA State 7)

Page 124 - First line of second paragraph should read:

The maximum number of DFA states is defined on line seven of Listing 39 to be ‘

Page 129 - First line should say “Listing 2.41”, not “Listing 2.40” Second line should say “line 119”, not “line 23”. The forth line should say “line 129”, not “line 31”. A replacement paragraph follows:

Several support functions are needed to do the work, all in Listing 2.41. The add_to_dstates() function on line 119 adds a new DFA state to the Dstates array and increments the number-of-states counter, Nstates. It returns the state number (the index in Dstates) of the newly added state. in_dstates() on line 139 of Listing 2.41 is passed a set of NFA states and returns the state number of an existing state that uses the same set, or -1 if there is no such state.

Page 129 - Replace lines 147 and 148 of Listing 2.41 with the following:

```
147     DFA_STATE *end = &Dstates[Nstates];  
148     for( p = Dstates ; p < end ; ++p )
```

Page 130 - Replace lines 193 and 194 of Listing 2.41 with the following

```
193     DFA_STATE *end = &Dstates[Nstates];  
194     for( p = Dstates; p < end ; ++p )
```

Page 139 - Replace line 198 of Listing 2.44 with the following:

```
198     SET    **end    = &Groups[Numgroups];
```

Replace Line 204 of Listing 139 with the following:

```
204     for( current = Groups; current < end; ++current )
```

Page 140 - Replace lines 229 and 230 of Listing 2.45 with the following

```
229     SET **end = &Groups[Numgroups];  
230     for( current = Groups; current < end; ++current )
```

Page 152 - Remove the `#include "date.h"` on line 3.

Page 157 - Replace line 117 with the following (I've added the `--argc` on the left):

```
117     for( ++argv, --argc; argc && *(p = *argv) == '-' ; ++argv, --argc )
```

Page 167 - 13th line from the bottom, “returns” should be “return.”

Page 167 - Third line from the bottom. Change the second comma to a semicolon.

Page 171 - Replace the first display (which now reads `noun→time|banana`) with the following:

`noun → fruit | banana`

Page 173 - Second display, text to the right of the `=L>` is missing. It should read as follows:

`compound_stmt =L> LEFT_CURLY stmt RIGHT_CURLY
=L> LEFT_CURLY RIGHT_CURLY`

Page 173 - First line after second display, change `expr→ε` to `stmt→ε`: The line should read:

The application of `stmt→ε` effectively removes the nonterminal from the derivation by ‘‘

Page 173 - 11th and 16th line from the bottom. Change `CLOSE_CURLY` to `RIGHT_CURLY`
`RIGHT_CURLY RIGHT_CURLY`

Page 175 - Figure 3.2. The line that reads ‘`4→DIGIT error`’ should read ‘`4→DIGIT 5.`’

5

Page 177 - Listing 3.2, Line 12, should read:

```
12     process_stmt( remember );
```

Page 178 - Third line below the one that starts with “Since” (next to the marginal note), replace “the the” with “the”: All three lines should read:

Since a `decl_list` can go to ϵ , the list could be empty. Parse trees for left and right recursive lists of this type are shown in Figure 3.5. Notice here, that the `decl_list→ε` Productions executed first or last.

production is the first list-element that's processed in the left-recursive list, and it's the last ‘‘

Page 179 - fourth line of second display, "declarator" should be "declarator_list". Line should read:

declarator_list → **TYPE** *declarator_list*

Page 180 - Table 3.2. Grammatical rules for “Separator Between List Elements, Zero elements okay” row are wrong. Replace the table with the following one:

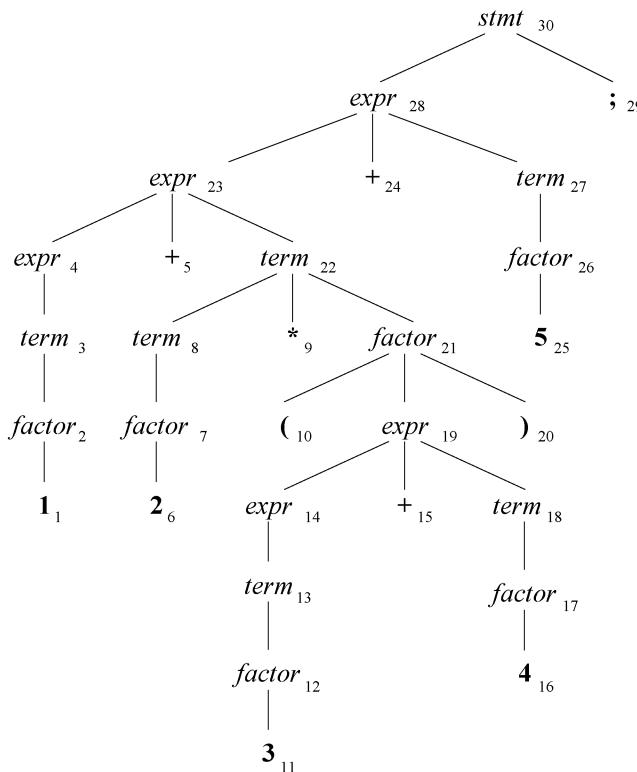
Table 3.2. List Grammars

	No Separator	
	Right associative	Left associative
At least one element	<i>list</i> → MEMBER <i>list</i> MEMBER	<i>list</i> → <i>list</i> MEMBER MEMBER
Zero elements okay	<i>list</i> → MEMBER <i>list</i> ϵ	<i>list</i> → <i>list</i> MEMBER ϵ
Separator Between List Elements		
	Right associative	Left associative
	<i>list</i> → MEMBER delim <i>list</i> MEMBER	<i>list</i> → <i>list</i> delim MEMBER MEMBER
At least one element	<i>opt_list</i> → <i>list</i> ϵ	<i>opt_list</i> → <i>list</i> ϵ
Zero elements okay	<i>list</i> → MEMBER delim <i>list</i> MEMBER	<i>list</i> → <i>list</i> delim MEMBER MEMBER

A MEMBER is a list element; it can be a terminal, a nonterminal, or a collection of terminals and nonterminals. If you want the list to be a list of terminated objects such as semicolon-terminated declarations, MEMBER should take the form: MEMBER → α **TERMINATOR**, where α is a collection of one or more terminal or nonterminal symbols.

Page 181 - Figure 3.6. Change all “statements” to “stmt” for consistency. Also change “expression” to “expr”. A new figure follows:

Figure 3.6. A Parse Tree for $1+2*(3+4)+5;$

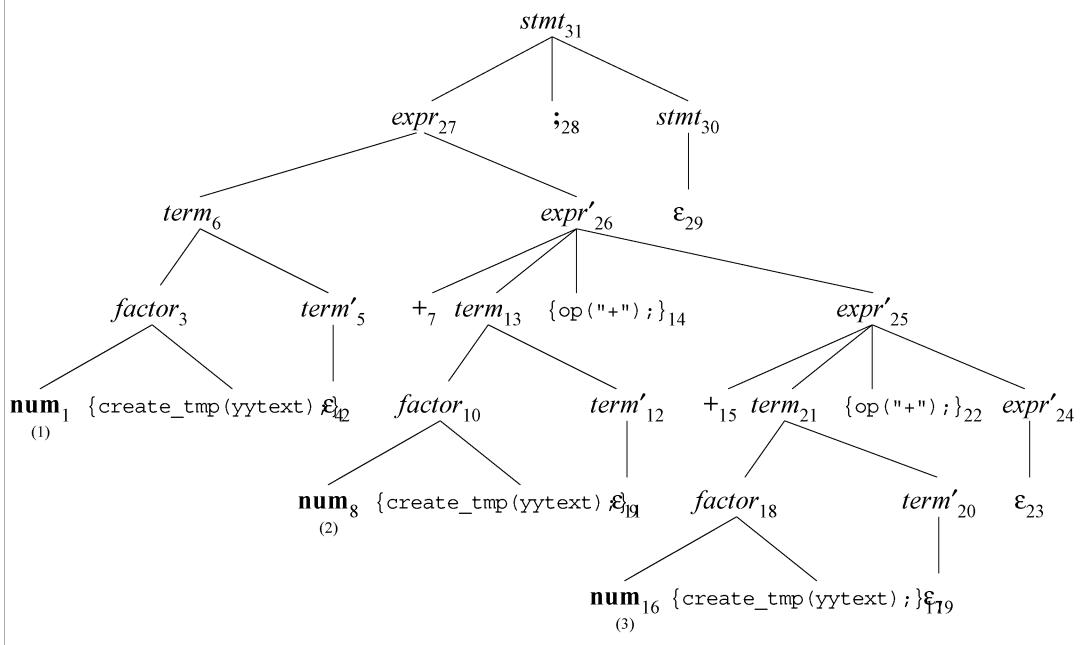


Page 183 - Display at bottom of page. Remove the exclamation point. The expression should read:

$$expr' \rightarrow + \ term \{ op ('+') ; \} \ expr'$$

Page 186 - Figure 3.8. Change caption to "Augmented parse tree for $1+2+3;$ " and change "statements" to "stmt". A new figure follows:

Figure 3.8. Augmented Parse Tree for $1+2+3$;



Page 188 - The caption for Listing 3.4 should say “Inherited” (not “Synthesized”) “Attributes.” A replacement caption follows:

Listing 3.4. *naive.c*—Code Generation with Inherited Attributes

Page 190 - First three lines beneath the figure (which start “right-hand side”) should read:

right-hand side of the production, as if they were used in the subroutine. For example, say that an attribute, t , represents a temporary variable name, and it is attached to an $expr$; it is represented like this in the subroutine representing the $expr$:

Page 196 - Figure 4.1. both instances of the word “number” should be in boldface.

number **number**

Page 208 - Third and eighth lines change “automata” to “automaton”. Replacement lines:

solution is a push-down automaton in which a state machine controls the activity on the ‘ ‘

The tables for the push-down automaton used by the parser are relatively straightfor-

Page 210 - Figure 4.6, `YY_d[factor]` [LP] should be "9," not "8."

9

Page 213 - First Paragraph, add **SEMICOLON** to the list on the third line. A replacement paragraph follows.

Production 1 is applied if a *statement* is on top of the stack and the input symbol is an **OPEN_CURLY**. Similarly, Production 2 is applied when a *statement* is on top of the stack and the input symbol is an **OPEN_PAREN**, **NUMBER**, **SEMICOLON**, or **IDENTIFIER** (an **OPEN_PAREN** because an *expression* can start with an **OPEN_PAREN** by Production 3, a **NUMBER** or **IDENTIFIER** because an *expression* can start with a *term*, which can, in turn, start with a **NUMBER** or **IDENTIFIER**). The situation is complicated when an *expression* is on top of the stack, however. You can use the same rules as before to figure out whether to apply Productions 3 or 4, but what about the ϵ production (Production 5)? The situation is resolved by looking at the symbols that can follow an *expression* in the grammar. If *expression* goes to ϵ , it effectively disappears from the current derivation (from the parse tree)—it becomes transparent. So, if an *expression* is on top of the stack, apply Production 5 if the current lookahead symbol can follow an *expression* (if it is a **CLOSE_CURLY**, **CLOSE_PAREN**, or **SEMICOLON**). In this last situation, there would be serious problems if **CLOSE_CURLY** could also start an *expression*. The grammar would not be LL(1) were this the case.

Page 213 - Last six lines should be replaced with the following seven lines.

ones. Initially, add those terminals that are at the far left of a right-hand side:

FIRST(stmt)	=	{ }
FIRST(expr)	=	{ ϵ }
FIRST(expr')	=	{PLUS, ϵ }
FIRST(term)	=	{ }
FIRST(term')	=	{TIMES, ϵ }
FIRST(factor)	=	{LEFT_PAREN, NUMBER}

Page 214 - Remove the first line beneath Table 4.14 [which starts ‘FIRST(factor)’].

Page 214 - Table 4.13, item (3), third line. Replace ‘is are’ with ‘are.’ A replacement table follows:

Table 3.3. Finding FIRST Sets

- (1) FIRST(**A**), where **A** is a terminal symbol, is {**A**}. If **A** is ϵ , then ϵ is put into the FIRST set.
- (2) Given a production of the form
 $s \rightarrow A \alpha$
where *s* is a nonterminal symbol, **A** is a terminal symbol, and α is a collection of zero or more terminals and nonterminals, **A** is a member of FIRST(*s*).
(3) Given a production of the form
 $s \rightarrow b \alpha$
where *s* and *b* are single nonterminal symbols, and α is a collection of terminals and nonterminals, everything in FIRST(*b*) is also in FIRST(*s*).
This rule can be generalized. Given a production of the form:
 $s \rightarrow \alpha B \beta$
where *s* is a nonterminal symbol, α is a collection of zero or more nullable nonterminals,[†] *B* is a single terminal or nonterminal symbol, and β is a collection of terminals and nonterminals, then FIRST(*s*) includes the union of FIRST(*B*) and FIRST(α). For example, if α consists of the three nullable nonterminals *x*, *y*, and *z*, then FIRST(*s*) includes all the members of FIRST(*x*), FIRST(*y*), and FIRST(*z*), along with everything in FIRST(*B*).

[†] A nonterminal is nullable if it can go to ϵ by some derivation. ϵ is always a member of a nullable nonterminal's FIRST set.

Page 214 - This is a change that comes under the “should be explained better” category and probably won’t make it into the book until the second edition. It confuses the issue a bit to put ϵ into the FIRST set, as per rule (1) in Table 4.13. (In fact, you could argue that it shouldn’t be there at all.) I’ve put ϵ into the FIRST sets because its presence makes it easier to see if a production is nullable (it is if ϵ is in the FIRST set). On the other hand, you don’t have to transfer the ϵ to the FOLLOW set when you apply the rules in Table 4.15 because ϵ serves no useful purpose in the FOLLOW set. Consequently, ϵ doesn’t appear in any of the FOLLOW sets that are derived on pages 215 and 216.

Page 214 - Bottom line and top of next page. Add ϵ to FIRST(expr), FIRST(expr'), and FIRST(term'). A replacement display, which replaces the bottom two lines of page 214 and the top four lines of page 215, follows:

FIRST(stmt)	=	{LEFT_PAREN, NUMBER, SEMI}
FIRST(expr)	=	{LEFT_PAREN, NUMBER, ϵ }
FIRST(expr')	=	{PLUS, ϵ }
FIRST(term)	=	{LEFT_PAREN, NUMBER}
FIRST(term')	=	{TIMES, ϵ }
FIRST(factor)	=	{LEFT_PAREN, NUMBER}

Page 216 - Add the following sentence to the end of item (2):

Note that, since ϵ serves no useful purpose in a FOLLOW set, it does not have to be transferred from the FIRST to the FOLLOW set when the current rule is applied. A replacement table follows:

Table 3.4. Finding FOLLOW Sets

- (1) If s is the goal symbol, eoi (the end-of-input marker) is in FOLLOW(s);
- (2) Given a production of the form:

$$s \rightarrow \dots a B \dots$$

where a is a nonterminal and B is either a terminal or nonterminal, FIRST(B) is in FOLLOW(a); To generalize further, given a production of the form:

$$s \rightarrow \dots a \alpha B \dots$$

where s and a are nonterminals, α is a collection of zero or more nullable nonterminals and B is either a terminal or nonterminal. FOLLOW(a) includes the union of FIRST(α) and FIRST(B). Note that, since ϵ serves no useful purpose in a FOLLOW set, it does not have to be transferred from the FIRST to the FOLLOW set when the current rule is applied.

- (3) Given a production of the form:

$$s \rightarrow \dots a$$

where a is the rightmost nonterminal on the right-hand side of a production, everything in FOLLOW(s) is also in FOLLOW(a). (I'll describe how this works in a moment.) To generalize further, given a production of the form:

$$s \rightarrow \dots a \alpha$$

where s and a are nonterminals, and α is a collection of zero or more nullable nonterminals, everything in FOLLOW(s) is also in FOLLOW(a).

Page 217 - Grammar in the middle of the page. Delete the pk and adm at the right edges of Productions 1 and 2.

Page 218 - Move the last line of page 217 to the top of the current page to eliminate the orphan.

Page 218 - Table 4.16, Replace the table with the following one (I've made several small changes). You may also want to move the widow at the top of the page to beneath the table while you're at it.

Table 4.16. Finding LL(1) Selection Sets

- (1) A production is *nullable* if the entire right-hand side can go to ϵ . This is the case, both when the right-hand side consists only of ϵ , and when all symbols on the right-hand side can go to ϵ by some derivation.
- (2) **For nonnullable productions:** Given a production of the form

$$s \rightarrow \alpha B \dots$$

where s is a nonterminal, α is a collection of one or more nullable nonterminals, and B is either a terminal or a nonnullable nonterminal (one that can't go to ϵ) followed by any number of additional symbols: the LL(1) select set for that production is the union of $\text{FIRST}(\alpha)$ and $\text{FIRST}(B)$. That is, it's the union of the FIRST sets for every nonterminal in α plus $\text{FIRST}(B)$. If α doesn't exist (there are no nullable nonterminals to the left of B), then $\text{SELECT}(s) = \text{FIRST}(B)$.

- (3) **For nullable productions:** Given a production of the form

$$s \rightarrow \alpha$$

where s is a nonterminal and α is a collection of zero or more nullable nonterminals (it can be ϵ): the LL(1) select set for that production is the union of $\text{FIRST}(\alpha)$ and $\text{FOLLOW}(s)$. In plain words: if a production is nullable, it can be transparent—it can disappear entirely in some derivation (be replaced by an empty string). Consequently, if the production is transparent, you have to look through it to the symbols that can follow it to determine whether it can be applied in a given situation.

Page 223 - Replace the last two lines on the page as follows:

Ambiguous productions, such as those that have more than one occurrence of a given nonterminal on their right-hand side, cause problems in a grammar because a unique parse tree is not generated for a given input. As we've seen, left factoring can be used to

Page 224 - The sentence on lines 13 and 14 (which starts with "If an ambiguous") should read as follows:

If an ambiguous right-hand side is one of several, then all of these right-hand sides must move as part of the substitution. For example, given:

Page 228 - 8th line from the bottom, the "Y" in "You" should be in lower case.

you can use a corner substitution to make the grammar self-recursive, replacing the 'Y' with a lowercase 'y'.

Page 222 - First line of text should read "Figures 4.5 and 4.6" rather than "Figure 4.6"
A replacement paragraph follows:

4.5 are identical in content to the ones pictured in Figures 4.5 and 4.6 on page 210. Note that the `yyd` table on lines 179 to 184 is not compressed because this output file was generated with the `-f` switch active. Were `-f` not specified, the tables would be pair compressed, as is described in Chapter Two. The `yy_act()` subroutine on lines 199 to 234 contains the switch that holds the action code. Note that `$` references have been translated to explicit value-stack references (`yy_vsp` is the value-stack pointer). The `YY_SYNCH` array on lines 243 to 248 is a -1 -terminated array of the synchronization tokens specified in the `%synch` directive.

Page 237 - The loop control on line 377 of Listing 4.6 won't work reliably in the 8086 medium or compact models. Replace lines 372–384 with the following:

```
372     int nterms;           /* # of terms in the production */
373     start = yy_pushtab[ production ];
374     for( end = start; *end; ++end )      /* After loop, end is positioned */
375     ;                                /* to right of last valid symbol */
376     count = sizeof(buf);
377     *buf = '\0';
378     for( nterms = end - start; --nterms >= 0 && count > 0 ; )      /* Assemble */
379     {
380         strncat( buf, yy_sym(--end), count );
381         if( (count -= strlen( yy_sym(end) + 1 )) < 1 )
382             break;
383         strncat( buf, " ", --count );
384     }
```

Page 242 - Pieces of the section heading for Section 4.9.2 are in the wrong font, and the last two lines are messed up. Replace with the following:

4.9.2 Occs and LLama Debugging Support—yydebug.c

This section discusses the debug-mode support routines used by the llama-generated parser in the previous section. The same routines are used by the occs-generated parser discussed in the next chapter. You should be familiar with the interface to the *curses*, window-management functions described in Appendix A before continuing.

Page 255 - Fourth line beneath the listing (starting with “teractive mode”), replace comma following the close parenthesis with a period. The line should read:

teractive mode (initiated with an *n* command). In this case, a speedometer readout that ‘‘

Page 271-303 - Odd numbered pages. Remove all tildes from the running heads.

Page 274 - Add the statement `looking_for_brace = 0;` between lines 179 and 180. Do it by replacing lines 180–190 with the following:

```
180                     looking_for_brace = 0;
181                 }
182             else
183             {
184                 if( c == '%' ) looking_for_brace = 1;
185                 else          output( "%c", c );
186             }
187         }
188     }
189     return CODE_BLOCK;
190 }
```

Page 278 - Fifth line from bottom. Replace `{create_tmp(yytext);}` with the following (to get the example to agree with Figure 4.9):

```
{rvalue(yytext);}
```

Page 282 - The last paragraph should read as follows (the Listing and Table numbers are wrong):

The recursive-descent parser for LLama is in Listing 4.25. It is a straightforward representation of the grammar in Table 4.19.

Page 312 - Listing 4.30. Replace lines 60 to 63 with the following:

```
60 PRIVATE int *Dtran; /* Internal representation of the parse table.  
61      * Initialization in make_yy_dtran() assumes  
62      * that it is an int [it calls memiset()].  
63      */
```

Page 315 - Listing 4.30. Replace lines 231 to 240 with the following:

```
231     nterms      = USED_TERMS + 1;           /* +1 for EOI */  
232     nnonterms   = USED_NONTERMS;  
233  
234     i = nterms * nnonterms;                 /* Number of cells in array */  
235  
236     if( !(Dtran = (int *) malloc(i * sizeof(*Dtran))) ) /* number of bytes */  
237         ferr("Out of memory\n");  
238  
239     memiset( Dtran, -1, i );                /* Initialize Dtran to all failures */  
240     ptab( Symtab, fill_row, NULL, 0 );       /* and fill nonfailure transitions. */
```

Page 330 - Listing 4.34, line 464. Delete everything except the line number.

Page 330 - Last two lines of second paragraph should read as follows:
bottom-up parse tables are created, below. Most practical LL(1) grammars are also LR(1) grammars, but not the other way around.

Page 330 - Add the right-hand side “| NUMBER” to the grammar in exercise 4.5. Also, align the table in Exercise 4.5 under the text so that it no longer extends into the gutter.

$$\begin{array}{lcl}
 \text{expr} & \rightarrow & -\text{expr} \\
 & | & * \text{expr} \\
 & | & \text{expr} * \text{expr} \\
 & | & \text{expr} / \text{expr} \\
 & | & \text{expr} = \text{expr} \\
 & | & \text{expr} + \text{expr} \\
 & | & \text{expr} - \text{expr} \\
 & | & (\text{expr}) \\
 & | & \text{NUMBER}
 \end{array}$$

Page 349 - Replace Table 5.5 with the following table:

Table 5.5. Error Recovery for $1++2$

Stack	Input	Comments
state: - parse: -	1 + + 2 -	Shift start state
state: 0 parse: \$	1 + + 2 -	Shift NUM (goto 1)
state: 0 1 parse: \$ NUM	+ + 2 -	Reduce by Production 3 ($\text{!T} \rightarrow \text{NUM}$) (Return to 0, goto 3)
state: 0 3 parse: \$!T	+ + 2 -	Reduce by Production 2 ($\text{!E} \rightarrow \text{!T}$) (Return to 0, goto 2)
state: 0 2 parse: \$!E	+ + 2 -	Shift !+ (goto 4)
state: 0 2 4 parse: \$!E !+	+ 2 -	ERROR (no transition in table) Pop one state from stack
state: 0 2 parse: \$!E	+ 2 -	There is a transition from 2 on !+ Error recovery is successful
state: 0 2 parse: \$!E	+ 2 -	Shift !+ (goto 4)
state: 0 2 4 parse: \$!E !+	2 -	Shift NUM (goto 1)
state: 0 2 4 1 parse: \$!E !+ NUM	2 -	Shift NUM (goto 1)
state: 0 2 4 1 parse: \$!E !+ NUM	-	Reduce by Production 3 ($\text{!T} \rightarrow \text{NUM}$) (Return to 4, goto 5)
state: 0 2 4 5 parse: \$!E !+ !T	-	Reduce by Production 1 ($\text{!E} \rightarrow \text{!E} \text{!+} \text{!T}$) (Return to 0, goto 2)
state: 0 2 parse: \$!E	-	Accept

Page 360 - Figure 5.6. The item immediately below the line in State 7 (ie. the first closure item) should be changed from $\text{!E} \rightarrow \text{. !E} \text{!*} \text{!F}$ to $\text{!T} \rightarrow \text{. !T} \text{!*} \text{!F}$

$\text{!T} \rightarrow \text{. !T} \text{!*} \text{!F}$

Page 361 - Third paragraph of section 5.6.2 (which starts “The FOLLOW”). Replace the paragraph with the following one:

The FOLLOW sets for our current grammar are in Table 5.6. Looking at the shift/reduce conflict in State 4, FOLLOW(!E) doesn’t contain a $!*$, so the SLR(1) method works in this case. Similarly, in State 3, FOLLOW(s) doesn’t contain a $!+$, so everything’s okay. And finally, in State 10, there is an outgoing edge labeled with a $!*$, but FOLLOW(!E) doesn’t contain a $!*$. Since the FOLLOW sets alone are enough to resolve the shift/reduce conflicts in all three states, this is indeed an SLR(1) grammar.

Page 361 - First paragraph in section 5.6.3 (which starts “Continuing our quest”). Replace the paragraph with the following one:

Many grammars are not as tractable as the current one—it’s likely that a FOLLOW set will contain symbols that also label an outgoing edge. A closer look at the machine yields an interesting fact that can be used to solve this difficulty. A nonterminal’s FOLLOW set includes *all* symbols that can follow that nonterminal in every possible context. The state machine, however, is more limited. You don’t really care which symbols can follow a nonterminal in every possible case; you care only about those symbols that can be in the input when you reduce by a production that has that nonterminal on its left-hand side. This set of relevant lookahead symbols is typically a subset of the complete FOLLOW set, and is called the *lookahead set*.

Some symbols in FOLLOW set are not needed.

Lookahead set.

Page 362 - Ninth line from bottom. Delete ‘only.’ A replacement for this, and the following four lines follows:

The process of creating an LR(1) state machine differs from that used to make an LR(0) machine only in that LR(1) items are created in the closure operation rather than LR(0) items. The initial item consists of the start production with the dot at the far left and $|$ as the lookahead character. In the grammar we’ve been using, it is:

Page 362 - Last line. Delete the period. The line should read: $x \rightarrow \gamma$

Page 363 - The C is in the wrong font in both the first marginal note and the first display (on the third line). It should be in Roman.

[$S \rightarrow \alpha \cdot x \beta, C$].

[$x \rightarrow !.\gamma, \text{FIRST}(\beta C)$].

Page 363 - Ninth line from bottom. Replace ‘new machine’ with ‘new states.’ A replacement paragraph follows:

The process continues in this manner until no more new LR(1) items can be created. The next states are created as before, adding edges for all symbols to the right of the dot and moving the dots in the kernel items of the new states. The entire LR(1) state machine for our grammar is shown in Figure 5.7. I’ve saved space in the Figure by merging together all items in a state that differ only in lookahead. The lookahead for all such items are shown on a single line in the right column of each state. Figure 5.8

shows how the other closure items in State 0 are derived. Derivations for items in States 2 and 14 of the machine are also shown.

Page 364 - Figure 5.7. About 3 inches from the left of the figure and $1\frac{3}{4}$ inches from the bottom, a line going from the box marked 2 to a circle with a B in it is currently labeled “*t e fNUM(.)*.” Delete the e.

Page 365 - The fourth line below Figure 5.7 should read:

best of both worlds. Examining the LR(1) machine in Figure 5.7, you are immediately ‘‘

Page 365 - Figure 5.7 (continued). The upper-case F in the second item of State 16 should be lower case.

$!T \rightarrow . !T !* !F$

Page 366 - First and third line under the Figure. The figure numbers are called out incorrectly in the text. The first three lines beneath the figure should read:

parenthesis first. The outer part of the machine (all of the left half of Figure 5.7 except States 6 and 9) handles unparenthesized expressions, and the inner part (States 6 and 9, and all of the right half of Figure 5.7) handles parenthesized subexpressions. The parser ‘‘

Page 370 - Listing 5.2, line 14. Change the sentence “Reduce by production *n*” to read as follows (leave the left part of the line intact):

Reduce by production *n*, *n* == -action.

Page 371 - Listing 5.2, line 16. Change the sentence “Shift to state *n*” to read as follows (leave the left part of the line intact):

Shift to state *n*, *n* == action.

Page 373 - Listing 5.4, line 6. Remove the yy in yylookahead. The corrected line looks like this:

```
6      do_this = yy_next( yy_action, state_at_top_of_stack(), lookahead );
```

Page 373 - Listing 5.4, line 29. Change rhs_len to rhs_length. The corrected line looks like this:

Page 373 - Last line. Change to read as follows (the state number is wrong):

shifts to State 1, where the only legal action is a reduce by Production 6 if the next input ‘‘

Page 374 - Paragraph beneath table, replace “you you” on third line with “you”. Entire replacement paragraph follows:

There’s one final caveat. You cannot eliminate a single-reduction state if there is a code-generation action attached to the associated production because the stack will have one fewer items on it than it should when the action is performed—you won’t be able to access the attributes correctly. In practice, this limitation is enough of a problem that occs doesn’t use the technique. In any event, the disambiguating rules discussed in the next section eliminate many of the single-reduction states because the productions that cause them are no longer necessary.

Page 387 - Listing 5.11, line 107. Add the word `short`. The repaired line looks like this:

```
107 #define YYF ((YY_TTYPE) ( unsigned short)~0 >>1 )
```

Page 388 - Second paragraph, third and fourth lines. Change “the largest positive integer” to “to the largest positive short `int`.” and remove the following “I’m.” The repaired lines read as follows:

subroutine, `yy_act_next()`, which I’ll discuss in a moment.) It evaluates to the largest positive short `int` (with two’s complement numbers). Breaking the macro down: ‘‘

Page 390 - Listing 5.12, line 199. Change the sentence “Reduce by production *n*” to read as follows (leave the left part of the line intact):

Reduce by production n, n == -action.

Page 390 - Listing 5.2, line 201. Change the sentence “Shift to state *n*” to read as follows (leave the left part of the line intact):

Shift to state n, n == action.

Page 397 - Replace line 209 of Listing 5.14 with the following:

```
209     YYD( yycomment ("Popping %s from state stack\n", tos); )
```

Page 398 - Listing 5.14, lines 219–222. Replace with the following code:

```
219     YY_vsp = YY_vstack + (YYMAXDEPTH - yystk_ele(Yy_stack)) ;
220 #ifdef YYDEBUG
221     yystk_p(Yy_dstack) = Yy_dstack +
222             (YYMAXDEPTH - yystk_ele(Yy_stack));
```

Page 403 - The loop control on Line 128 of Listing 5.15 doesn't work reliably in the 8086 compact or large model. To fix it, replace Line 97 of Listing 5.15 (p. 403) with the following (and also see change for next page):

```
97     int i ;
```

Page 404 - Replace line 128 of Listing 5.15 with the following:

```
128     for(i = (pp - prod->rhs) + 1; --i >= 0; --pp)
```

Page 425 - Lines 585–594. Replace with the following code:

```
585     if( nclose )
586     {
587         assort( closure_items, nclose, sizeof(ITEM*), item_cmp );
588         nitems = move_eps( cur_state, closure_items, nclose );
589         p      = closure_items + nitems;
590         nclose -= nitems;
591
592         if( Verbose > 1 )
593             pclosure( cur_state, p, nclose );
594     }
```

Page 440 - Listing 5.33, replace the code on lines 1435 to 1438 with the following (be sure that the quote marks on the left remain aligned with previous lines):

```

1435      "    action < 0    -- Reduce by production n,  n == -action.",
1436      "    action == 0   -- Accept. (ie. Reduce by production 0.)",
1437      "    action > 0    -- Shift to state n,  n == action.",
1438      "    action == YYF  -- error.",

```

Page 447 - Line 14. Change "hardly every maps" to "hardly ever maps". The line should read:

ideal machine hardly ever maps to a real machine in an efficient way, so the generated '
'

Page 452 - First line below Listing 6.2. Change "2048" to "1024". Line should read:

In addition to the register set, there is a 1024-element, 32-bit wide stack, and two '
'

Page 453 - Figure 6.1. Around the middle of the figure. Change "2048" to "1024". Line should read:

1,024
32-bit
lwords

Page 466 - Listing 6.11, lines 53 and 54, change sp to __sp (two underscores).

```

53 #define push(n)          ( __sp ) ->l = ( lword ) (n)
54 #define pop(t)           (t) ( __sp++ ) ->l )

```

Page 471 - Ninth line. Replace with fp+4 with fp+16. Replace the last four lines of the paragraph with the following lines:

call() subroutine modifies wild, it just modifies the memory location at fp+16, and on the incorrect stack, ends up modifying the return address of the calling function. This means that call() could work correctly, as could the calling function, but the program would blow up when the calling function returned.

Page 475 - Listing 6.19, Line 80. Replace the first (b) with a (s). The line should now read:

```

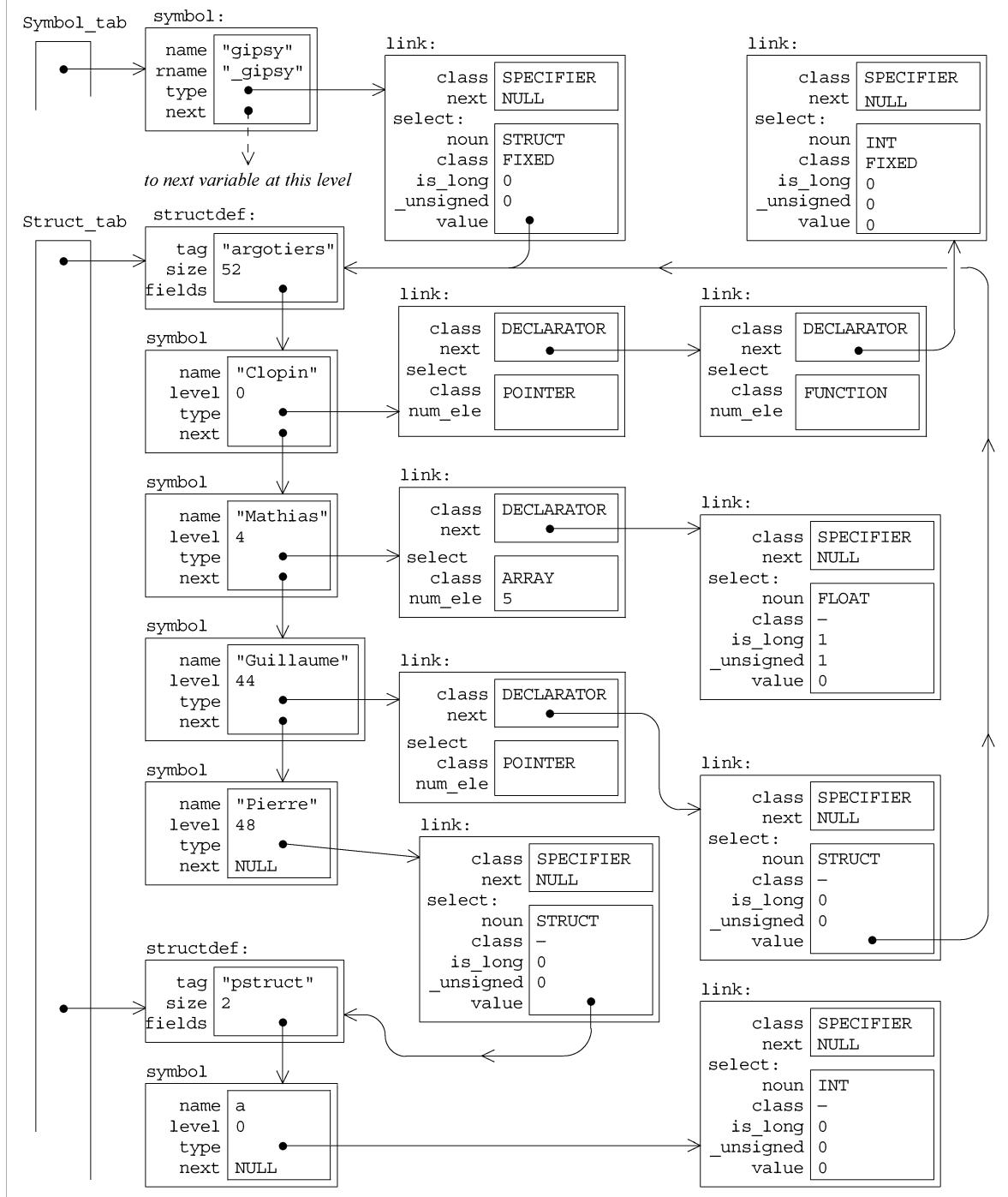
80 #define BIT(b,s)     if( (s) & (1 << (b)) )

```

Page 486 - Replace last line of first paragraph (which now reads '6.10') with the following:
6.11.

```
129 #define IS_SPECIFIER(p) ((p) && (p)->class==SPECIFIER )
130 #define IS_DECLARATOR(p) ((p) && (p)->class==DECLARATOR )
131 #define IS_ARRAY(p) ((p) && (p)->class==DECLARATOR && (p)->DCL_TYPE==ARRAY )
132 #define IS_POINTER(p) ((p) && (p)->class==DECLARATOR && (p)->DCL_TYPE==POINTER )
133 #define IS_FUNCT(p) ((p) && (p)->class==DECLARATOR && (p)->DCL_TYPE==FUNCTION)
134 #define IS_STRUCT(p) ((p) && (p)->class==SPECIFIER && (p)->NOUN == STRUCTURE )
135 #define IS_LABEL(p) ((p) && (p)->class==SPECIFIER && (p)->NOUN == LABEL )
136
137 #define IS_CHAR(p) ((p) && (p)->class == SPECIFIER && (p)->NOUN == CHAR )
138 #define IS_INT(p) ((p) && (p)->class == SPECIFIER && (p)->NOUN == INT )
139 #define IS_UINT(p) ( IS_INT(p) && (p)->UNSIGNED )
140 #define IS_LONG(p) ( IS_INT(p) && (p)->LONG )
141 #define IS ULONG(p) ( IS_INT(p) && (p)->LONG && (p)->UNSIGNED )
142 #define IS_UNSIGNED(p) ((p) && (p)->UNSIGNED )
```

Figure 6.13. Representing a Structure in the Symbol Table



Page 500 - First sentence below figure should start “The subroutine”

The subroutine in Listing 6.29 manipulates declarators: `add_declarator()` adds

Page 503 - Listing 6.31, line 281, should read as follows:

281

```
(p1->DCL_TYPE==ARRAY && (p1->NUM_ELE != p2->NUM_ELE)) )
```

Page 520 - Replace line 71 of Listing 6.37 with the following line:

71 ('.') | ('\\.') | ('\\{o}({o}{o}?){}') | ('\\x{h}({h}{h}?){}') |

Replace line 76 of Listing 6.37 with the following line:

76 ({d}+|{d}+\.{d}*|{d}*\.{d}+) ([eE][\+-]?{d}+)?[fF]? return FCON ;

Page 521 - Listing 6.37, Lines 138–143. Replace as follows:

```
138 typedef struct /* Routines to recognize keywords. A table */
139 { /* lookup is used for this purpose in order to */
140     char *name; /* minimize the number of states in the FSM. A */
141     int val; /* KWORD is a single table entry. */
142 }
143 KWORD;
```

Page 524 - Second line of last paragraph, remove period after **TYPE** and change “Listing 6.38” to Listing 6.39. The repaired line should read as follows:

reduces `typeSpecifier→TYPE` (on line 229 of Listing 6.39). The associated action ‘

Page 527 - First line below Listing 6.40. Change “Listing 6.40” to Listing 6.39. The repaired line should read as follows:

216 to 217 of Listing 6.39.) There are currently three attributes of interest: `$1` and `$$` ‘

Page 553 - Listing 6.56, line 10. Change L0 to (L0*4).

```
10 #define T(n) (fp- (L0*4) - (n*4))
```

Page 556 - Listing 6.58, line 539. Change %s to (%s*4)

```
539     yycode( "#define T(n) (fp- (%s*4) - (n*4))\n\n", Vspace );
```

Page 558 - Listing 6.60, lines 578–582. Replace with the following:

```
578     discard_link_chain(existing->type); /* Replace existing type      */
579     existing->type    = sym->type;   /* chain with the current one.*/
580     existing->etype   = sym->etype;
581     sym->type    = sym->etype = NULL; /* Must be NULL for discard_- */
582     }                      /* symbol() call, below.      */
```

Page 558 - Listing 6.60, lines 606 and 607. i is not used and the initial offset should be 8. Replace lines 606 and 607 with the following:

```
606     int offset = 8;           /* First parameter is always at BP(fp+8):      */
607     /* 4 for the old fp, 4 for the return address. */
```

Page 560 - Page 560, Listing 6.61. Replace lines 578–580 with the following:

```
578     : LC
579             {
580                 if( ++Nest_lev == 1 )
581                     loc_reset(); }
```

Page 573 - Fifth line from the bottom. Insert a period after “needed”. The line should read:

needed. The stack is shrunk with matching additions when the variable is no longer ‘‘

Page 574 - Figure 6.18, the L0 in the `#define T(n)` should be (L0*4).

```
#define T(n) (fp- (L0*4) - (n*4))
```

Page 578 - First paragraph. There's an incomplete sentence on the first line. Replace with the following paragraph and add the marginal note:

The cell is marked as “in use” on line 73. The Region element corresponding to the first cell of the allocated space is set to the number of stack elements that are being allocated. If more than one stack element is required for the temporary, adjacent cells that are part of the temporary are filled with a place marker. Other subroutines in Listing 6.64 de-allocate a temporary variable by resetting the equivalent Region elements to zero, de-allocate all temporary variables currently in use, and provide access to the high-water mark. You should take a moment and review them now.

Marking a stack cell as “in use”.

Page 590 - Listing 6.69, line 214. Replace with the following:

```
214      case CHAR:      return BYTE_PREFIX;
```

Page 598 - Third paragraph, second line (which starts “type int for”), replace line with the following one:

type int for the undeclared identifier (on line 62 of Listing 6.72). The ‘

Page 601 - First line beneath Listing 6.76. Replace “generate” with “generated”:

So far, none of the operators have generated code. With Listing 6.77, we move ‘

Page 607 - Listing 6.84, line 328: Change the || to an &&.

```
328      if( !IS_INT(offset->type) && !IS_CHAR(offset->type) )
```

Page 608 - Fonts are wrong in all three marginal notes. Replace them with the ones given here.

Operand to * or [] must be array or pointer.

Attribute synthesized by * and [] operators.

Rules for forming lvalues and rvalues when processing * and [] .

Page 613 - First line of last paragraph is garbled. Since the fix affects the entire paragraph, an entire replacement paragraph follows. Everything that doesn't fit on page 613 should be put at the top of the next page.

call()
unary→NAME

The `call()` subroutine at the top of Listing 6.87 generates both the `call` instruction and the code that handles return values and stack clean up. It also takes care of implicit subroutine declarations on lines 513 to 526. The action in `unary→NAME` creates a symbol of type `int` for an undeclared identifier, and this symbol eventually ends up here as the incoming attribute. The `call()` subroutine changes the type to “function returning `int`” by adding another link to the head of the type chain. It also clears the `implicit` bit to indicate that the symbol is a legal implicit declaration rather than an undeclared variable. Finally, a C-code `extern` statement is generated for the function.

Page 617 - Listing 6.87, line 543. Replace `nargs` with `nargs * SWIDTH`.

543	<code>gen("+=%s%d" , "sp", nargs * SWIDTH); /* sp is a byte pointer, */</code>
-----	--

Page 619 - Listing 6.88, line 690. Delete the `->name`. The repaired line should look like this:

690	<code>gen("EQ" , rvalue(\$1), "0");</code>
-----	--

Disk only. *Page 619, Listing 6.88. Added semantic-error checking to first (test) clause in `?:` operator. Tests to see if its an integral type. Insert the following between lines 689 and 690:*

<code>if(!IS_INT(\$1->type)) yyerror("Test in ?: must be integral\n");</code>
--

Page 619 - Listing 6.88. Replace line 709 with the following line:

709	<code>gen("=" , \$\$->name , rvalue(\$7));</code>
-----	--

Page 644 - Lines 895 and 896. *There is a missing double-quote mark on line 895, insertion of which also affects the formatting on line 896. Replace lines 895 and 896 with the following:*

```
895             gen("goto%s%d", L_BODY,      $5 );
896             gen(":%s%d",     L_INCREMENT, $5 );
```

Page 648 - Listing 6.107, line 1, change the 128 to 256.

```
1 #define CASE_MAX 256          /* Maximum number of cases in a switch */
```

Page 649 - Listing 6.108. Add the following two lines between lines 950 and 951: (These lines will not have numbers on them, align the first p in pop with the g in gen_stab... on the previous line.)

```
pop( S_brk );
pop( S_brk_label );
```

Page 658 - First paragraph of section 7.2.1 should be replaced with the following one:

A *strength reduction* replaces an operation with a more efficient operation or series of operations that yield the same result in fewer machine clock cycles. For example, multiplication by a power of two can be replaced by a left shift, which executes faster on most machines. ($x*8$ can be done with $x<<3$.) You can divide a positive number by a power of two with a right shift ($x/8$ is $x>>3$ if x is positive) and do a modulus division by a power of two with a bitwise AND ($x \% 8$ is $x\&7$).

Page 671 - Figure 7.1, third subfigure from the bottom. In initial printings, the asterisk that should be at the apex of the tree had dropped down about $\frac{1}{2}$ inch. Move it up in these printings. In later printings, there are two asterisks. Delete the bottom one.

Page 681 - Listing A.1, lines 19 and 20 are missing semicolons. Change them to the following:

```
19     typedef long      time_t; /* for the VAX, may have to change this */
20     typedef unsigned size_t; /* for the VAX, may have to change this */
```

Page 682 - Listing A.1, line 52. Delete all the text on the line, but leave the asterisk at the far left.

Page 688 - 6th line from the bottom. Remove “though” at start of line.

it might introduce an unnecessary conversion if the stack type is an int, short, or char. These multiple type conversions will also cause portability problems if the

`stack_err()` macro evaluates to something that won't fit into a `long` (like a `double`).

Page 690 - Last line, “calling conventions” should not be hyphenated.

Page 696 - Listing A.4, line 40. The comment is wrong. The line should read as follows:

```
40 #define _DIFFERENCE 2 /* (x in s1) and (x not in s2) */
```

Page 702 - Listing A.6. Change lines 98–102 to the following:

```
98     /* Enlarge the set to "need" words, filling in the extra words with zeros.
99      * Print an error message and exit if there's not enough memory.
100     * Since this routine calls malloc, it's rather slow and should be
101     * avoided if possible.
102    */
```

Page 706 - Listing A.8, line 330. Change `unsigned` to `int`:

```
330     int ssize; /* Number of words in src set */
```

Page 713 - Third line from the bottom. “`nextsym()`” should be in the Courier font. Replace the last three lines on the page with the following:

passing the pointer returned from `find_sym()` to `nextsym()`, which returns either a pointer to the next object or `NULL` if there are no such objects. Use it like this:

Page 719 - Second line above Listing A.17, change “`tree`” to “`table`”:

cial cases. The `delsym()` function, which removes an arbitrary node from the table, is shown in Listing A.17.

Page 722 - Listing A.19, line 221. Replace with:

```
221     return (*User_cmp)( (void*) (*p1 + 1), (void*) (*p2 + 1) );
```

Page 729 - Listing A.26, line 50. Delete the `(two required)` at the end of the line.

Page 736 - Listing A.33, line 34. Change to the following:

```
34 PUBLIC void stop_prnt() {}
```

Page 737 - Listing A.33, line 97. Change to the following:

```
97 char *str, *fmt, *argp;
```

Page 739 - The swap statement in the code in the middle of the page is incorrect.

Here is a replacement display:

```
int array[ ASIZE ];
int i, j, temp ;

for( i = 1; i < ASIZE; ++i )
    for( j = i-1; j >= 0; --j )
        if( array[j] > array[j+1] )
            swap( array[j], array[j+1] );
```

Page 743 - Listing A.36. Delete the text (but not the line number) on line 4 (which now says `#include <fcntl.h>`).

Page 745 - Change caption of Listing A.38 to the following:

Listing A.38. *memiset.c*— Initialize Array of `int` to Arbitrary Value

Page 755 - Third line from the bottom. Delete the exclamation point. The line should read:

images (25×80×2=4,000 bytes for the whole screen), and that much memory may not be

Page 758 - Listing A.45, line 49. Remove the semicolon at the far right of the line.

Page 768 - Listing A.61. Remove the exclamation point from the caption.

Page 776 - Line above heading for section A.11.2.2. Delete the `void`.

Page 797 - Replace Lines 50–59 of Listing A.84 with the following:

```
50     case '\b':  if( buf > sbuf )
51         {
52             --buf;    wprintf( win, " \b" );
53         }
54     else
55     {
56         wprintf( win, " " );
57         putchar('\007');
58     }
59 break;
```

Page 803 - Just above heading for section B.2. The line should read:

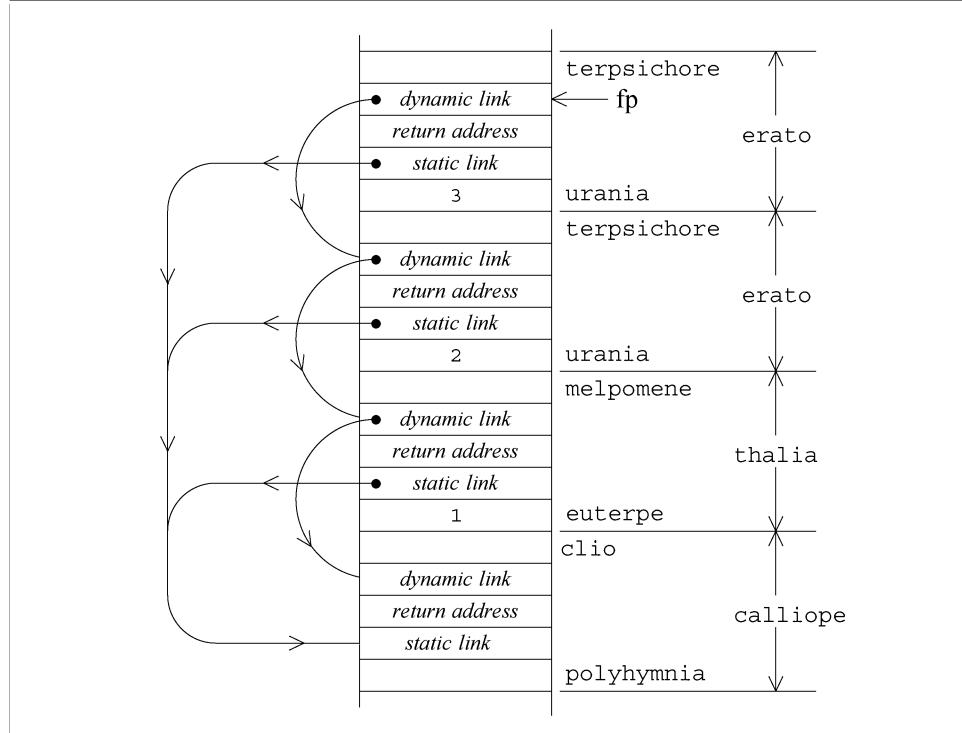
set to 100; otherwise, arg is set to 1.

Page 803 - Replace the last paragraph on the page with the following one:

The stack, when the recursive call to erato is active, is shown in Figure B.2. There is one major difference between these stack frames and C stack frames: the introduction of a second pointer called the *static link*. The *dynamic link* is the old frame pointer, just as in C. The static link points, not at the previously active subroutine, but at the parent subroutine in the nesting sequence—in the declaration. Since erato and thalia are both nested inside calliope, their static links point at calliope’s stack frame. You can chase down the static links to access the local variables in the outer routine’s ‘‘

Page 804 - Replace Figure B.2 with the following figure.

Figure B.2. Pascal Stack Frames



Page 805 - Top of the page. Replace the top seven lines (everything up to the paragraph that begins “This organization”) with the following (you’ll have to move the rest of the text on page 805 down to make room):

stack frame. For example, `clio` can be accessed from `erato` with the following C-code:

```
r0.pp = WP(fp + 4); /* r0 = static link */
_x    = W(r0.pp - 8); /* x = clio */
```

You can access `polyhymnia` from `erato` with:

```
r0.pp = WP(fp + 4); /* r0 = static link */
_x    = W(r0.pp + 8); /* x = clio */
```

Though it’s not shown this way in the current example, it’s convenient for the frame pointer to point at the static, rather than the dynamic link to make the foregoing indirection a little easier to do. The static links can be set up as follows: Assign to each subroutine a *declaration level*, equivalent to the nesting level at which the subroutine is declared. Here, `calliope` is a level 0 subroutine, `erato` is a level 1 subroutine, and so forth. Then:

- **If a subroutine calls a subroutine at the same level**, the static link of the called subroutine is identical to the static link of the calling subroutine.
- **If a subroutine at level N calls a subroutine at level $N+1$** , the static link of the called subroutine points at the static link of the calling subroutine.
- **If a subroutine calls a subroutine at a lower (more outer) level**, use the following algorithm:

```
i = the difference in levels between the two subroutines;
p = the static link in the calling subroutine's stack frame;
while( -i >= 0 )
    p = *p;
the static link of the called subroutine = p;
```

Note that the difference in levels (i) can be figured at compile time, but you must chase down the static links at run time. Since the static link must be initialized by the calling subroutine (the called subroutine doesn’t know who called it), it is placed beneath the return address in the stack frame.

Page 806 - Change caption and title of Listing C.1 as follows:

Listing C.1. A Summary of the C Grammar in Chapter Six.

The ^ and \$ metacharacters.

Page 819 - First full paragraph. Replace the “the the” on the fourth line with a single “the.”

A replacement paragraph follows:

The ^ and \$ metacharacters work properly in all MS-DOS input modes, regardless of whether lines end with `\r\n` or a single `\n`. Note that the newline is not part of the lexeme, even though it must be present for the associated expression to be recognized. Use `and\r\n` to put the end of line characters into the lexeme. (The `\r` is not required in UNIX applications, in fact it’s an error under UNIX.) Note that, unlike the vi editor `^$` does not match a blank line. You’ll have to use an explicit search such as `\r\n\r\n` to find empty lines.

Page 821 - Listing D.1, replace lines 14 and 15 with the following:

```
14     while( yylex() )  
15     ;
```

Page 821 - First two lines beneath the listing. Delete both lines and replace with the following text:

LeX and `yylen` is adjusted accordingly. Zero is returned at end of file, `-1` if the lexeme is too long.¹³

Page 821 - Replace Footnote 13 at the bottom of the page with the one at the bottom of the page you are now reading.

Page 828 - Listing D.5.; in order to support the `uL` suffix, replace line 16 with the following:

```
16  suffix  ([UuLl] | [uU] [lL]) /* Suffix in integral numeric constant */
```

Page 841 - Replace the code on the first five lines of Listing E.2 with the following five lines:

```
1 %term ID    /* an identifier */  
2 %term NUM   /* a number */  
3 %left PLUS  /* + */  
4 %left STAR   /* * */  
5 %left LP RP  /* ( and ) */
```

Page 843 - Paragraph starting `-c[N],`, 2nd line. Delete “e” in “switche.”

Page 860 - 15 lines from the bottom. Remove the underscores. The line should read:

is from stack picture six to seven. `t0` and `t1` were put onto the stack when the rvalues ‘‘

¹³ UNIX `lex` doesn’t return `-1` and it doesn’t modify the `yytext` or `yylen`; it just returns the next input character.

Page 861 - *Figure E.5. Remove all underscores. The figure should be replaced with the following one:*

Figure E.5. A Parse of A*2



Page 862 - *Move caption for Listing E.9 to the left. (It should be flush with the left edge of the box.)*

Page 871 - *Figure E.6. Label the line between States 5 and 7 with STAR.*

STAR

Page 880 - *Listing E.17, Lines 84 and 85, replace with the following:*

```
84     yycode("public word t0, t1, t2, t3;\n");
85     yycode("public word t4, t5, t6, t7;\n");
```

Page 886 - *Listing E.19. Replace lines 51 to 57 with the following:*

```
51 { 0} 512, line 42 : {$1=$2=newname(); }
52 { 1} 513, line 42 : {freename($0); }
53 { 2} 514, line 48 : {$1=$2=newname(); }
54 { 3} 515, line 49 : { yycode("%s+=%s\n", $$,$0); freename($0); }
55 { 4} 516, line 56 : {$1=$2=newname(); }
56 { 5} 517, line 57 : { yycode("%s*=%s\n", $$,$0); freename($0); }
57 { 6} 518, line 61 : { yycode("%s=%0.*s\n", $$,yylen,yytext); }
```

Page 887 - *Listing E.19. Replace lines 46 and 54 with the following:*

```
46 { yycode("%s+=%s\n", $$,$0); freename($0); } expr'
54 { yycode("%s*=%s\n", $$,$0); freename($0); } term'
```

Page 888 - *Listing E.19. Replace line 58 with the following:*

```
58 factor : NUM_OR_ID { yycode("%s=%0.*s\n", $$, yyleng, yytext); }
```

Disk only. I made several changes to *searchen.c* (p. 747) to make the returned path names more consistent (everything's now mapped to a UNIX-style name). Also added a disk identifier when running under DOS.

Disk only. Insert the following into the brace-processing code, between lines 115 and 116 of *parser.lex* on page 273:

```
if( i == '\n' && in_string )
{
    lerror( WARNING,
            "Newline in string, inserting \"\n\"");
    in_string = 0;
}
```

Disk only. The *do_unop* subroutine on page 604 (Line 177 of Listing 6.79) wasn't handling incoming constant values correctly and it wasn't doing any semantic-error checking at all. It's been replaced by the following code. (Instructions are now generated only if the incoming value isn't a constant, otherwise the constant value at the end of the link chain is just modified by performing the indicated operation at compile time.)

```
177 value *do_unop( op, val )
178 int op;
179 value *val;
180 {
181     char *op_buf = "=?" ;
182     int i;
183
184     if( op != '!' ) /* ~ or unary - */
185     {
186         if( !IS_CHAR(val->type) && !IS_INT(val->type) )
187             yyerror( "Unary operator requires integral argument\n" );
188
189         else if( IS_UNSIGNED(val->type) && op == '-' )
190             yyerror( "Minus has no meaning on an unsigned operand\n" );
191
192         else if( IS_CONSTANT( val->type ) )
193             do Unary_Const( op, val );
194         else
195         {
196             op_buf[1] = op;
197             gen( op_buf, val->name, val->name );
198         }
199     }
200     else /* ! */
201     {
202         if( IS_AGGREGATE( val->type ) )
203             yyerror("May not apply ! operator to aggregate type\n");
204     }
}
```

Listing 5.10. continued...

```
205     else if( IS_INT_CONSTANT( val->type ) )
206         do_unary_const( '!', val );
207     else
208     {
209         gen( "EQ",           rvalue(val), "0"                  ); /* EQ(x, 0)      */
210         gen( "goto%s%d",   L_TRUE,          i = tf_label() ); /* goto T000;    */
211         val = gen_false_true( i, val );           /* fall thru to F */
212     }
213 }
214 return val;
215 }
216 /-----*/
217 do_unary_const( op, val )
218 int op;
219 value *val;
220 {
221     /* Handle unary constants by modifying the constant's value. */
222
223     link *t = val->type;
224
225     if( IS_INT(t) )
226     {
227         switch( op )
228         {
229             case '^': t->V_INT = ~t->V_INT; break;
230             case '-': t->V_INT = -t->V_INT; break;
231             case '!': t->V_INT = !t->V_INT; break;
232         }
233     }
234     else if( IS_LONG(t) )
235     {
236         switch( op )
237         {
238             case '^': t->V_LONG = ~t->V_LONG; break;
239             case '-': t->V_LONG = -t->V_LONG; break;
240             case '!': t->V_LONG = !t->V_LONG; break;
241         }
242     }
243     else
244         yyerror("INTERNAL do_unary_const: unexpected type\n");
245 }
```

Disk only. Page 506, Listing 6.32, lines 453–457: Modified subroutine `type_str()` in file `symtab.c` to print the value of an integer constant. Replaced the following code:

```
if( link_p->NOUN != STRUCTURE )
    continue;
else
    i = sprintf(buf, " %s", link_p->V_STRUCT->tag ?
                link_p->V_STRUCT->tag : "untagged");
```

with this:

```

if( link_p->NOUN == STRUCTURE )
    i = sprintf(buf, " %s", link_p->V_STRUCT->tag ?
                link_p->V_STRUCT->tag : "un>tagged");

else if(IS_INT(link_p) ) sprintf(buf, "%d", link_p->V_INT );
else if(IS_UINT(link_p) ) sprintf(buf, "%u", link_p->V_UINT );
else if(IS_LONG(link_p) ) sprintf(buf, "%ld", link_p->V_LONG );
else if(IS ULONG(link_p) ) sprintf(buf, "%lu", link_p->V ULONG );
else continue;

```

Disk only. *Page 241, Listing 4.7, lines 586–589 and line 596, and page 399 Listing 5.14, lines 320–323 and line 331.* The code in *llama.par* was not testing correctly for a NULL return value from *ii_ptext()*. The problem has been fixed on the disk, but won't be fixed in the book until the second edition. The fix looks like this:

```

if( yytext = (char *) ii_ptext() ) /* replaces llama.par lines 586-589 */
{
    /* and occs.par, lines 320-323 */
    yylineno = ii_plineno();
    tchar = yytext[ yylen = ii_plength() ];
    yytext[yylen] = '\0';
}
else /* no previous token */
{
    yytext = "";
    yylen = yylineno = 0;
}

if( yylineno ) /* replaces llama.par, line 596 */
    ii_ptext()[ ii_plength() ] = tchar; /* and occs.par, line 331 */ */

```

Disk only. The *ii_look()* routine (in Listing 2.7 on page 47) doesn't work in the 8086 large or compact models. The following is ugly, but it works everywhere:

```

1 int ii_look( n )
2 {
3     /* Return the nth character of lookahead, EOF if you try to look past
4      * end of file, or 0 if you try to look past either end of the buffer.
5      * We have to jump through hoops here to satisfy the ANSI restriction
6      * that a pointer can not go to the left of an array or more than one
7      * cell past the right of an array. If we don't satisfy this restriction,
8      * then the code won't work in the 8086 large or compact models. In
9      * the small model---or in any machine without a segmented address
10     * space, you could do a simple comparison to test for overflow:
11     *     uchar *p = Next + n;
12     *     if( !(Start_buf <= p && p < End_buf) )
13     *         overflow
14     */
15
16     if( n > (End_buf-Next) ) /* (End_buf-Next) is the # of unread */
17     return Eof_read ? EOF : 0 ; /* chars in the buffer (including */
18     /* the one pointed to by Next). */
19
20     /* The current lookahead character is at Next[0]. The last character */
21     /* read is at Next[-1]. The --n in the following if statement adjusts */
22     /* n so that Next[n] will reference the correct character. */
23

```

Listing 5.11. continued...

```
24     if( --n < -(Next-Start_buf) )      /* (Next-Start) is the # of buffered */
25         return 0;                      /* characters that have been read. */
26
27     return Next[n];
28 }
```

