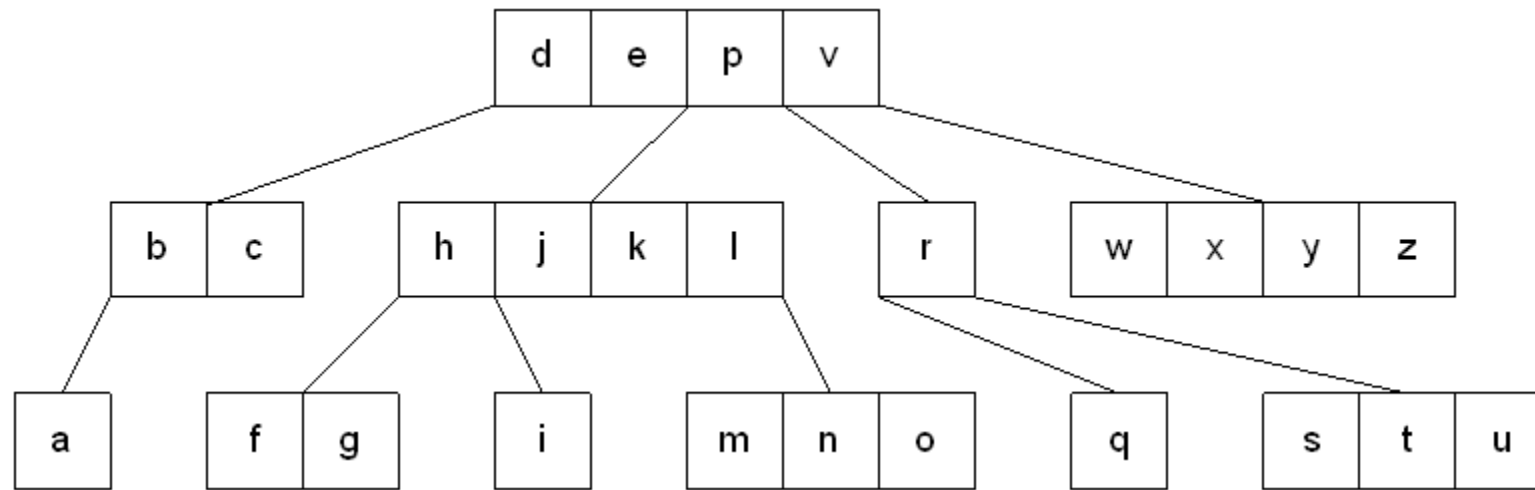


# M-way tree

Module-2

# M-way search tree:

- The **m-way** search trees are multi-way trees which are generalized versions of binary trees where each node contains multiple elements.
- In an m-way tree of order **m**, each node contains a maximum of **m – 1** elements and **m** children.
- Multiway tree of order 5



- Elements are present in ascending order at every level.

# M-way search tree:

- The goal of m-way search tree of height  $h$  calls for the search time complexity is  $O(h)$ .
- An m-way search tree of  $n$  elements ranges from a minimum height of  $\log_m (n+1)$  to a maximum of  $n$ .

Ex: Suppose you want to build a 3-way search tree ( $m = 3$ ) with  $n=26$  elements.

According to the formula:  $h_{\min} = \log_m (n+1) = \log_3 (26+1) = \log_3 (27) = 3$

- The number of elements in an m-way search tree of height  $h$  ranges from a minimum of  $h$  to a maximum of  $m^h - 1$ .

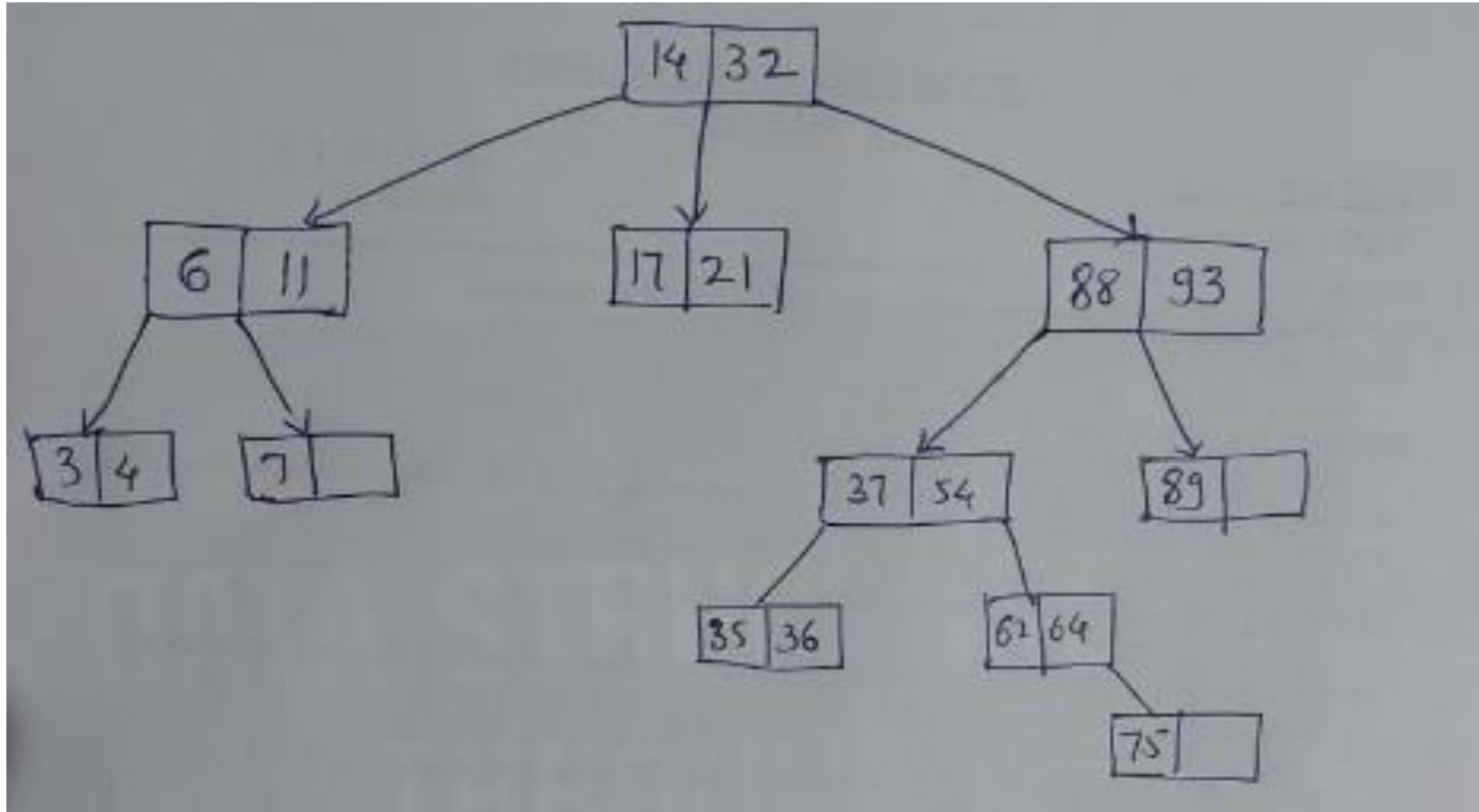
# Insertion in m-way tree

Eg. Create a m-way search tree with order=3.

Insert elements 14, 32, 11, 6, 7, 4, 3, 88, 93, 54, 37, 21, 17, 89, 62, 64, 75, 35, 36

- 1) Insert keys in node till key count ==  $m-1$
- 2) If Node full(no of keys ==  $m-1$ ) insert key to left sub tree node where key  $k_{new} < key_{parent}$
- 3) If Node full(no of keys ==  $m-1$ ) insert key to right sub tree node where key  $k_{new} > key_{parent}$
- 4) Repeat Process

# Insertion in m-way tree



# Deletion in m-way tree

Case 1: Deletion of key with no subtrees --> Simply delete key from node

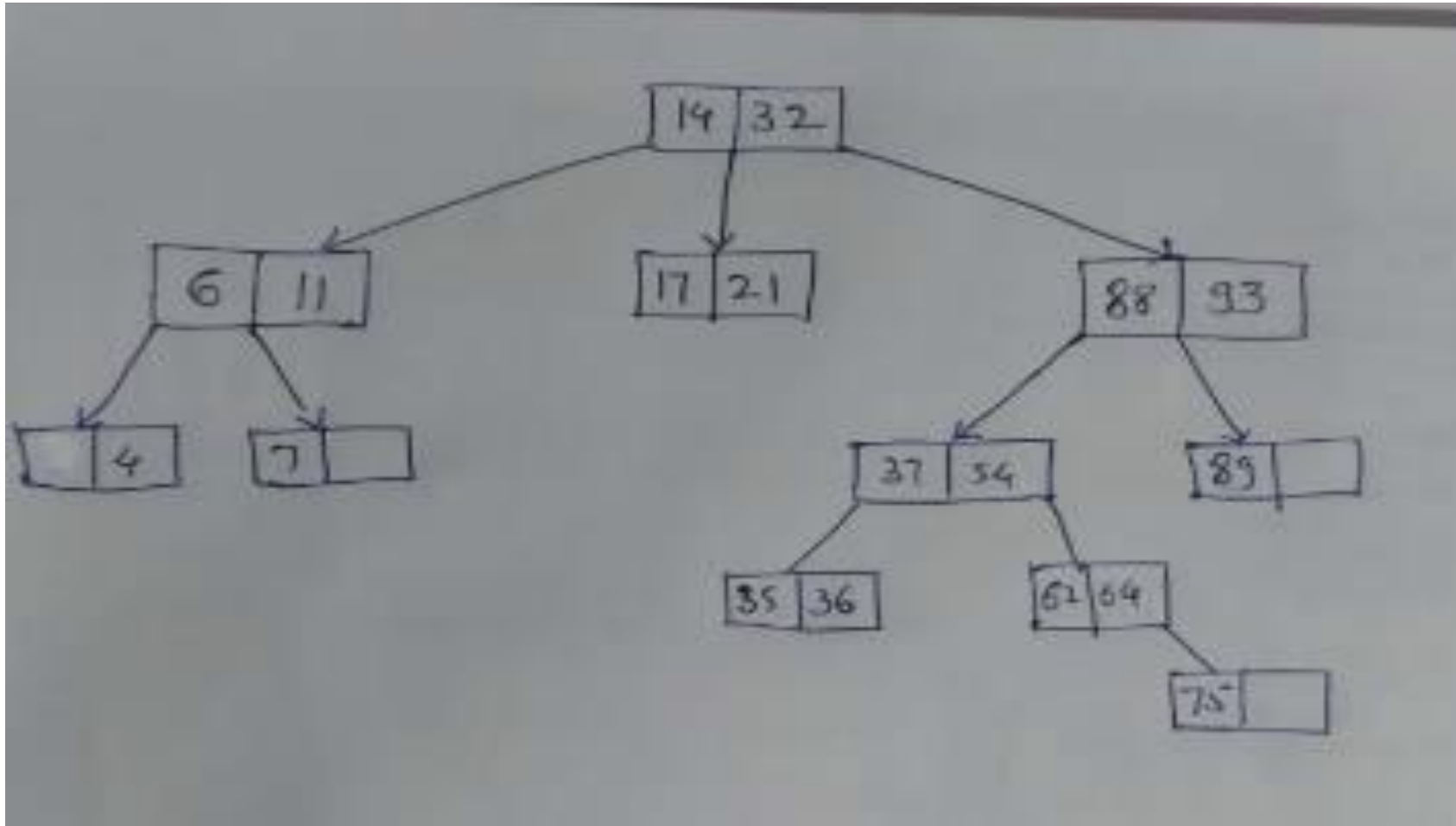
Case 2a: Deletion of key with left subtree --> Replace largest value from left subtree with value to be deleted. Delete largest value from left subtree.

Case 2b: Deletion of key with right subtree --> Replace smallest value from right subtree with value to be deleted. Delete smallest value from right subtree.

Case 3: Deletion of key with both left and right subtrees --> Replace smallest value for RST or Replace largest value from LST with the key value to be deleted. Then delete this replaced value.

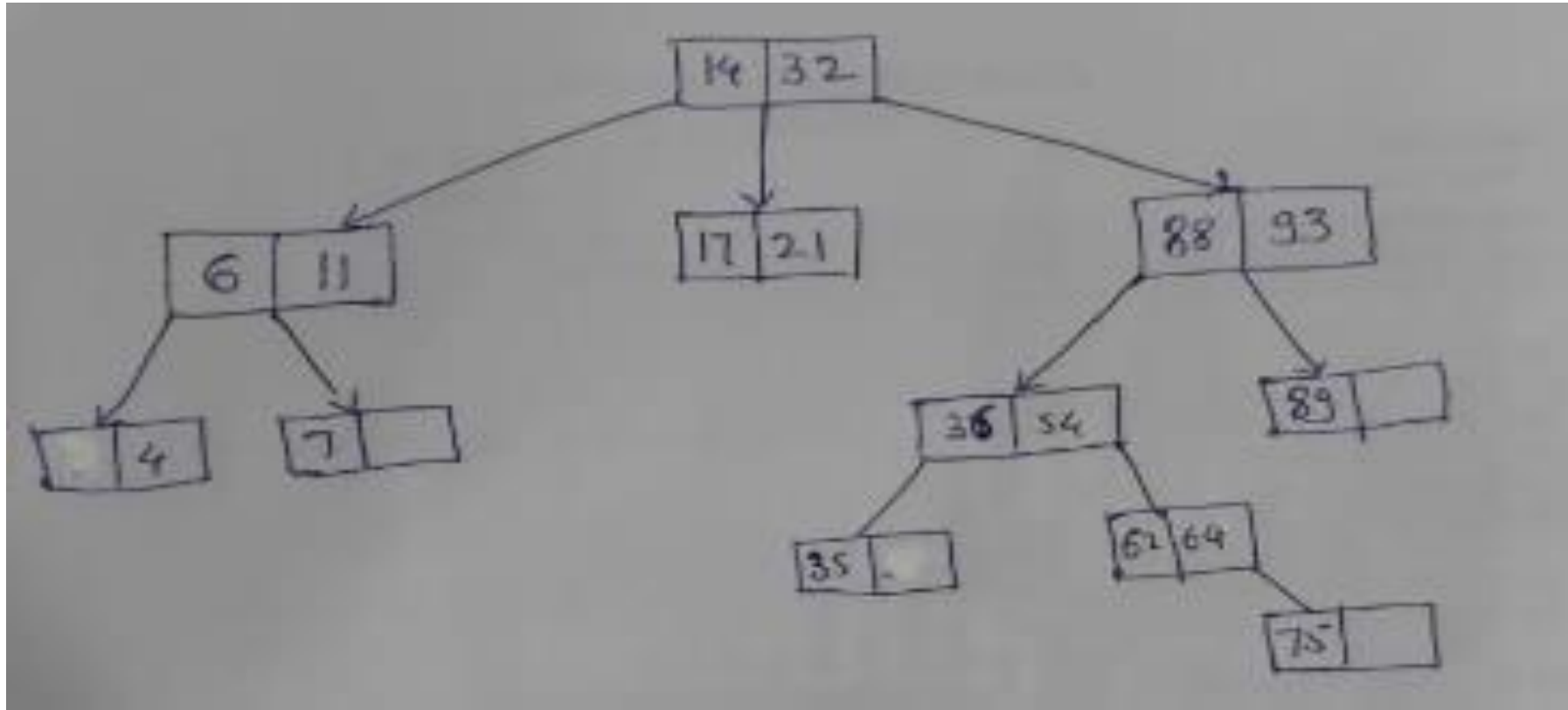
# Deletion in m-way tree

Delete 3 (case1)



# Deletion in m-way tree

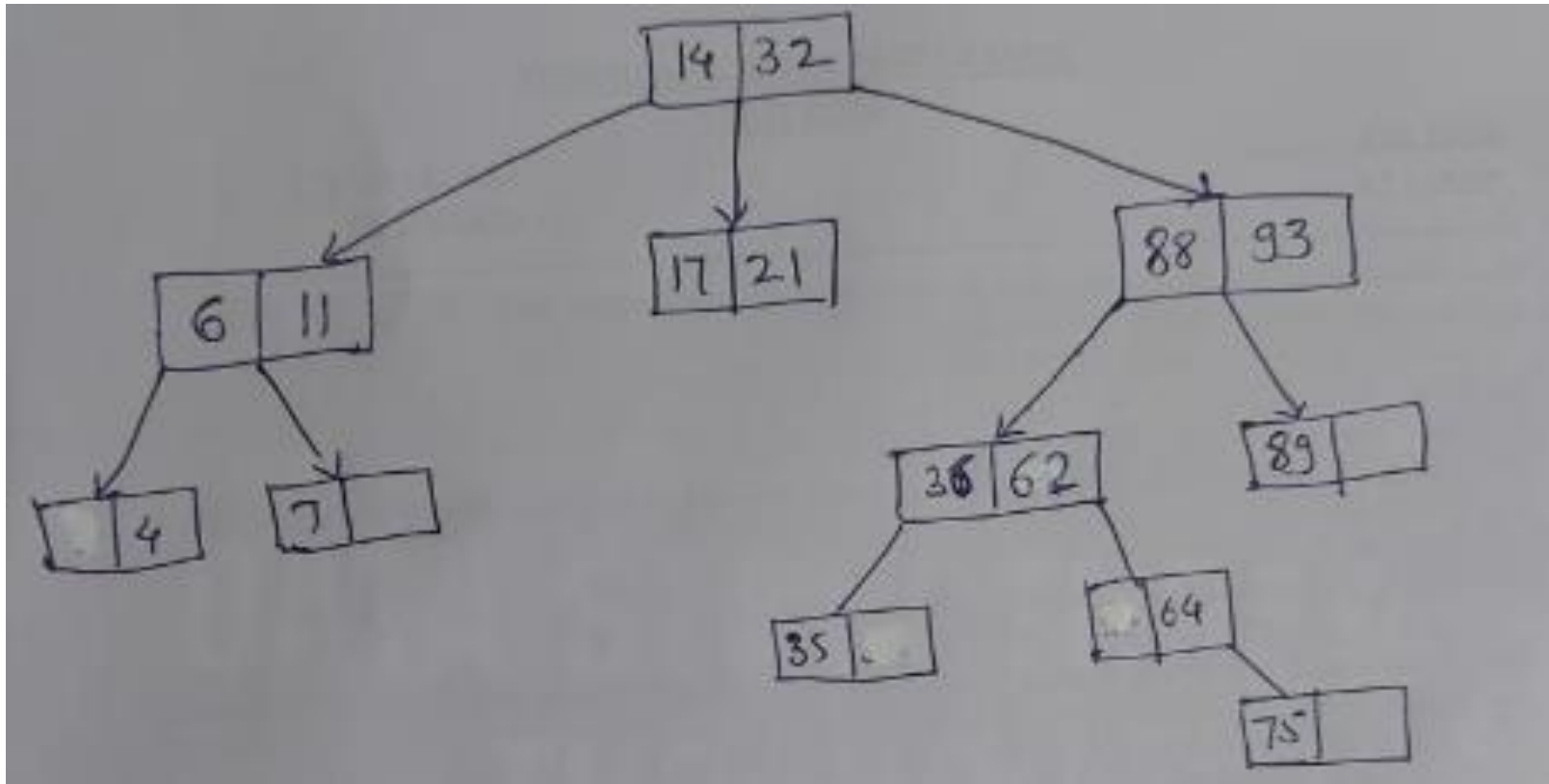
Delete 37 (case 2a)





# Deletion in m-way tree

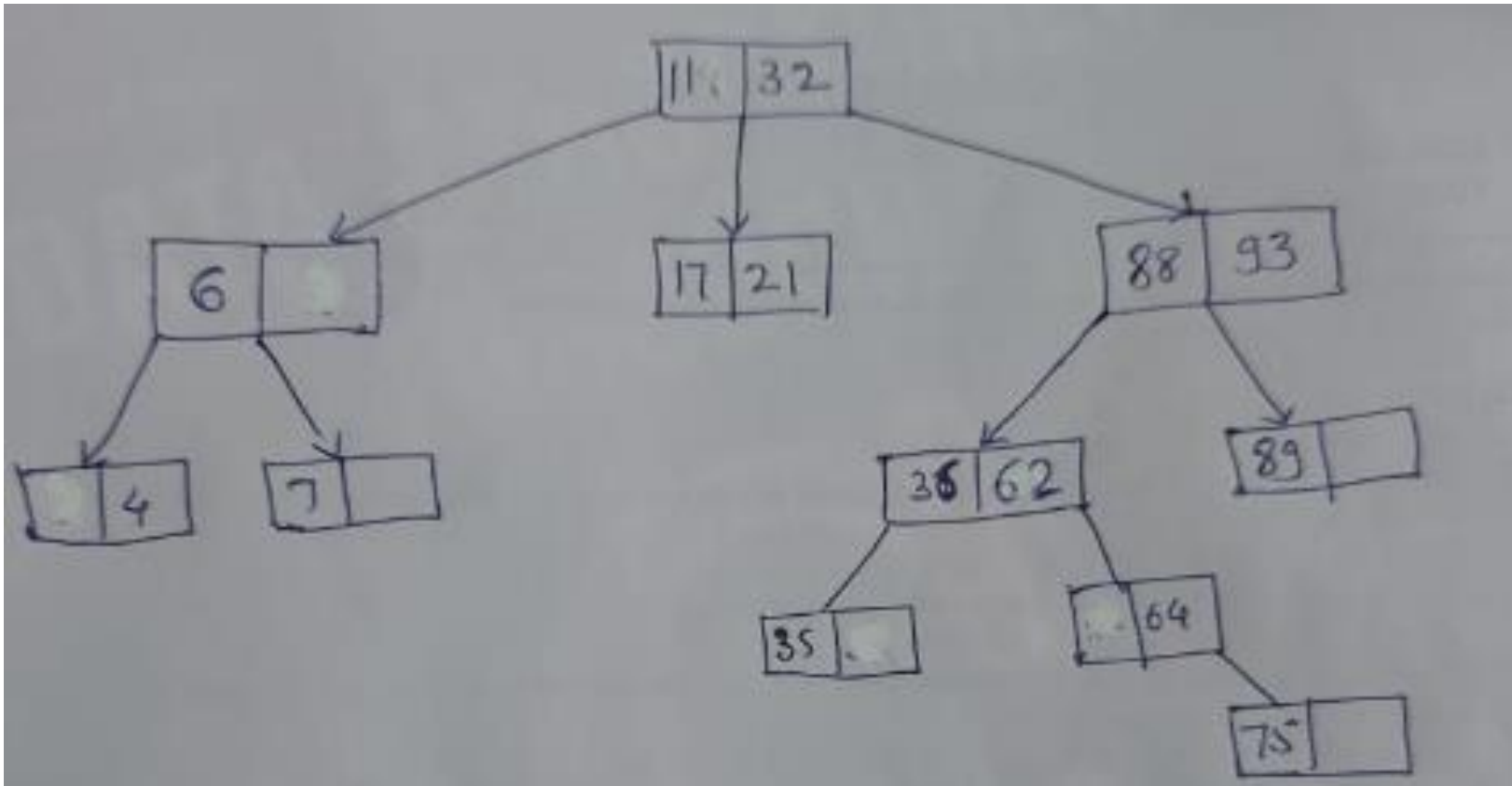
Delete 54 (case 2b)



# Deletion in m-way tree

Delete 14 (case 3)

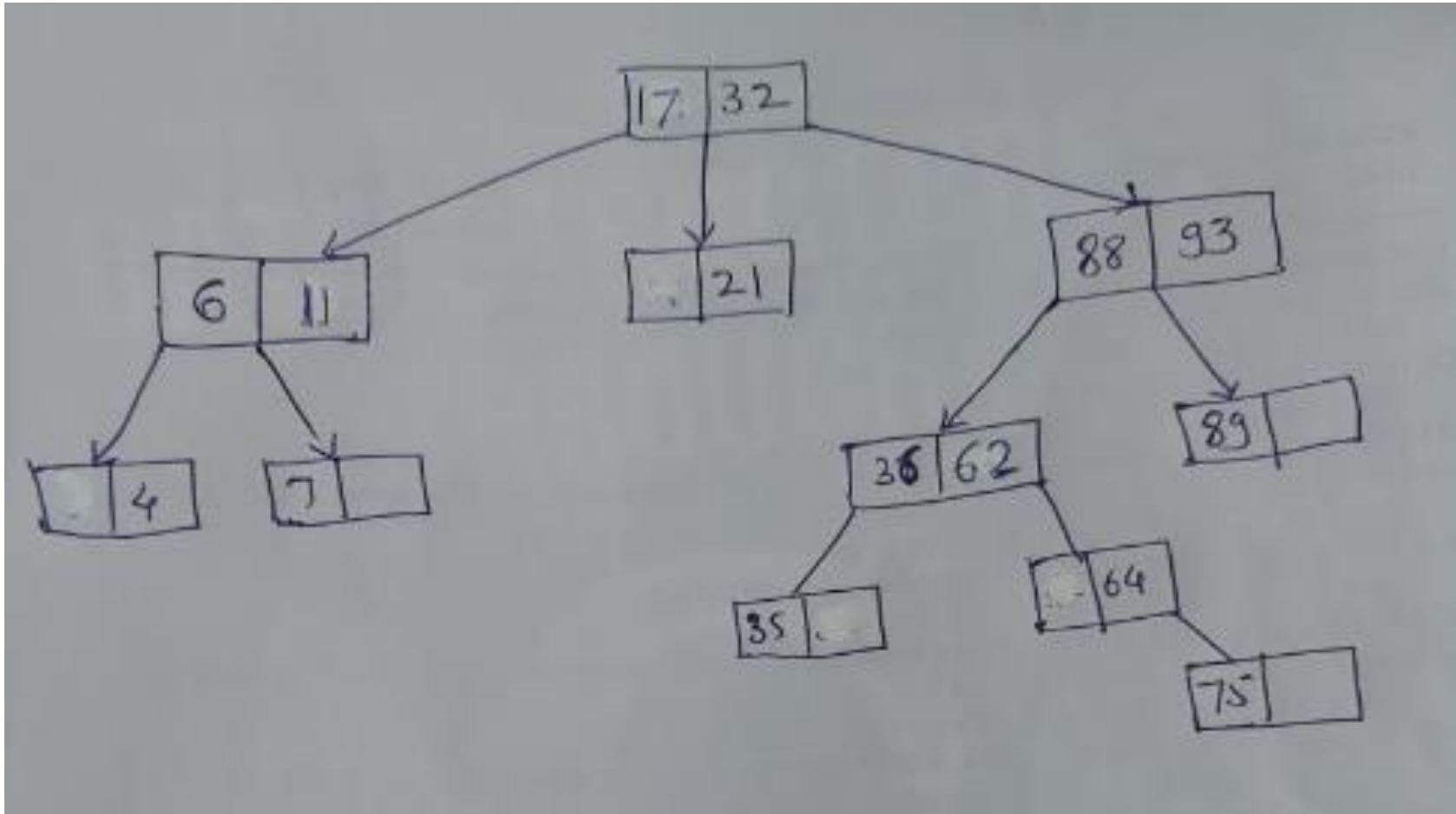
1<sup>st</sup> way to replace with largest value from left subtree



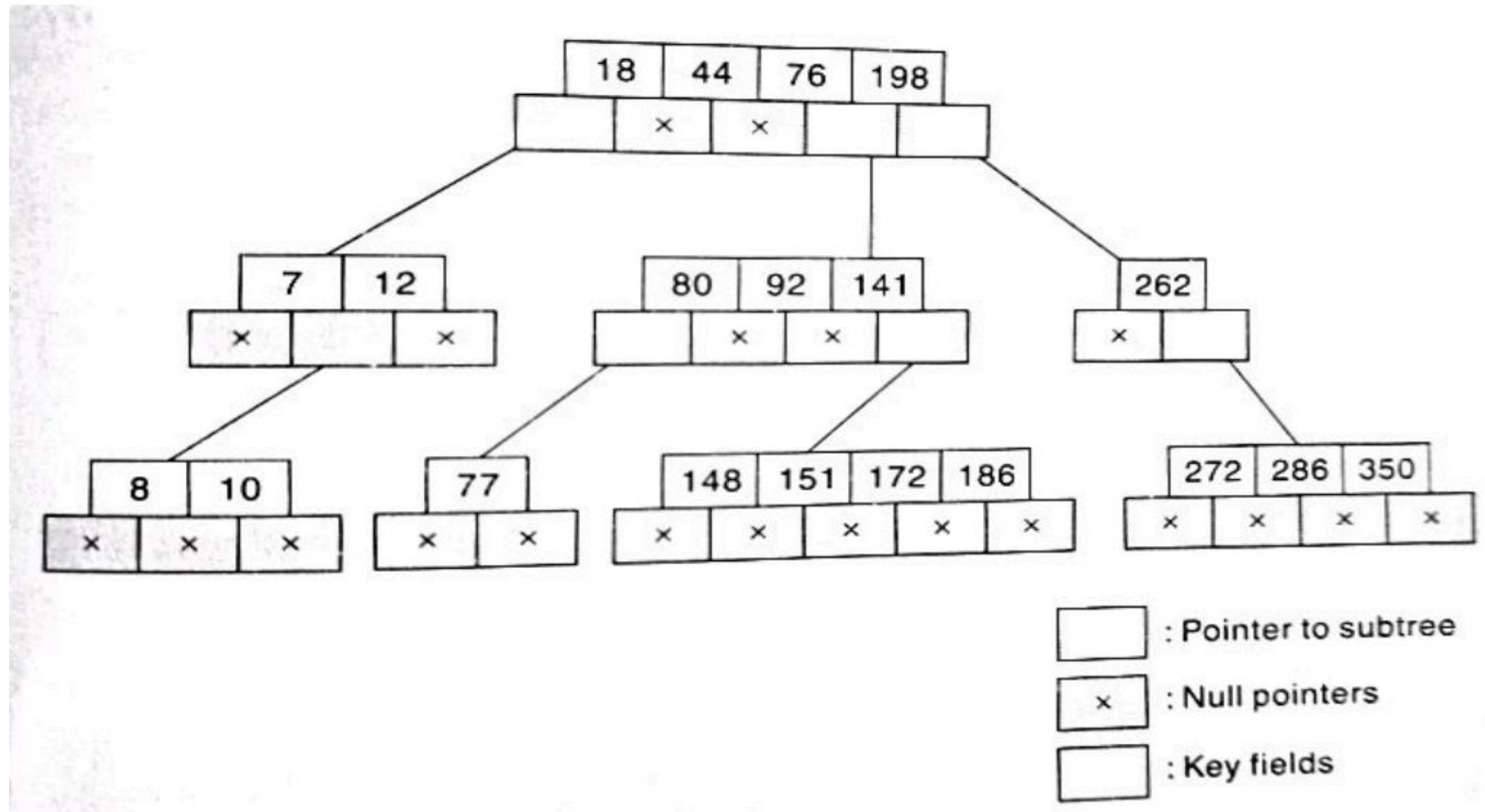
# Deletion in m-way tree

Delete 14 (case 3)

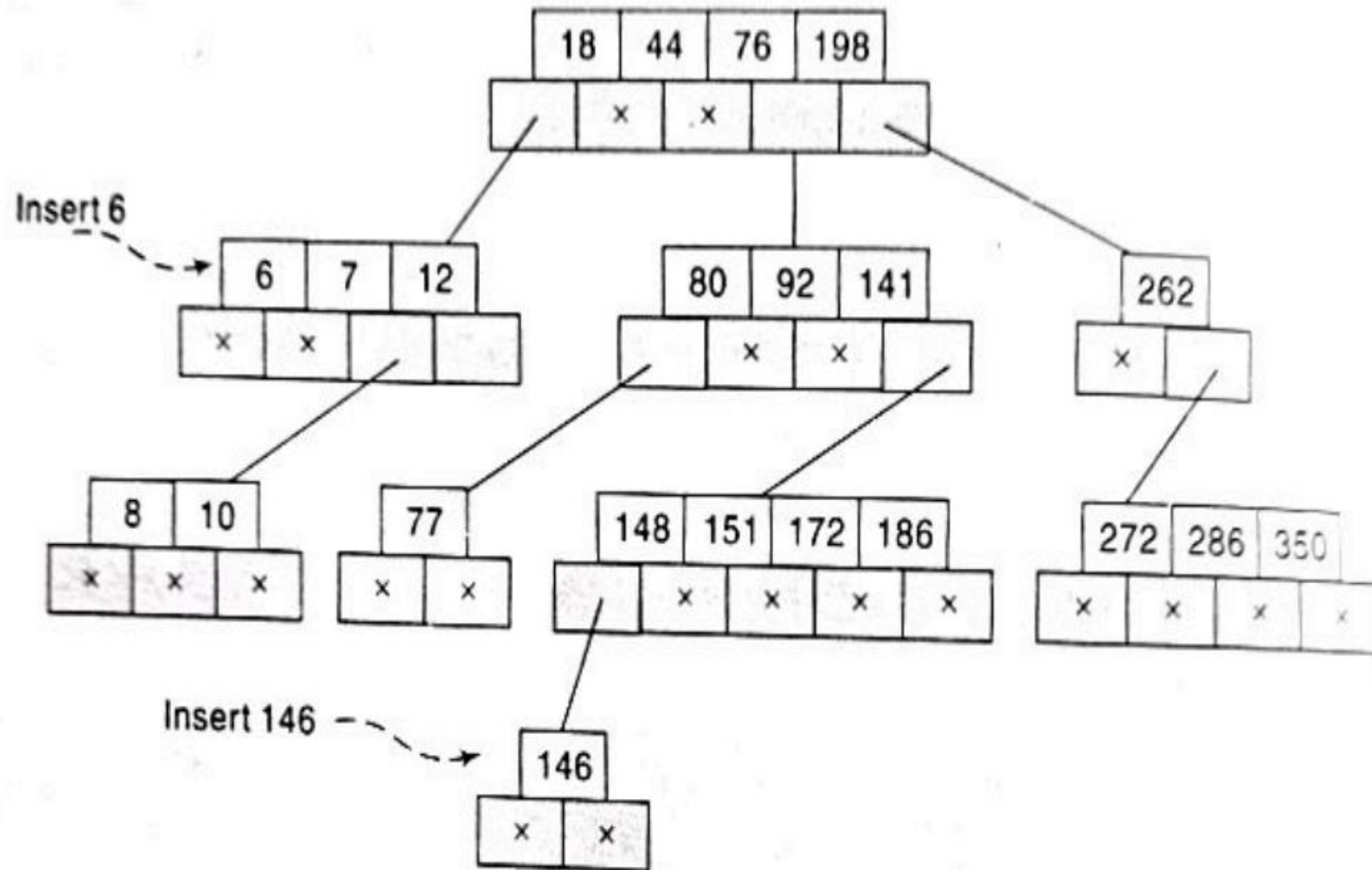
2<sup>nd</sup> way to replace with smallest value from right subtree



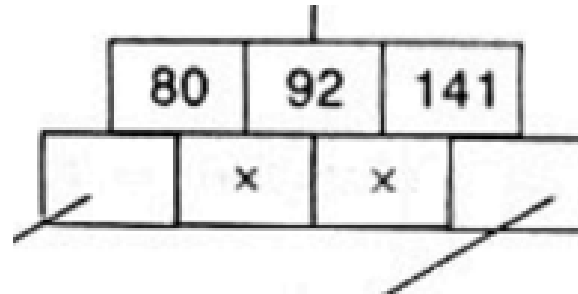
Given a following m-way tree. Insert 6 and 146 in above tree.



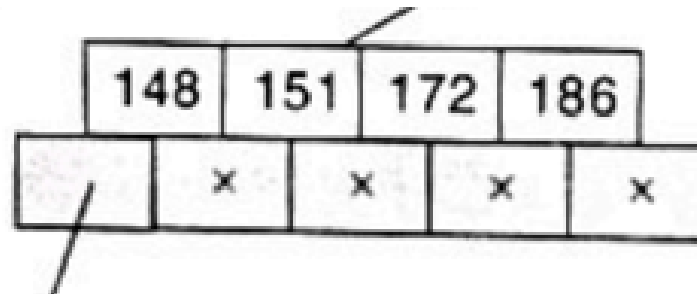
Given a following m-way tree. Insert 6 and 146 in above tree.



- Here we can not insert 146 in following node.



- Because if 146 will be added after 141 and this will lead to violation of property of m-way tree. In that case the following node will become left child of 146.



- Therefore we will add 146 as left child of above node(of key=148)