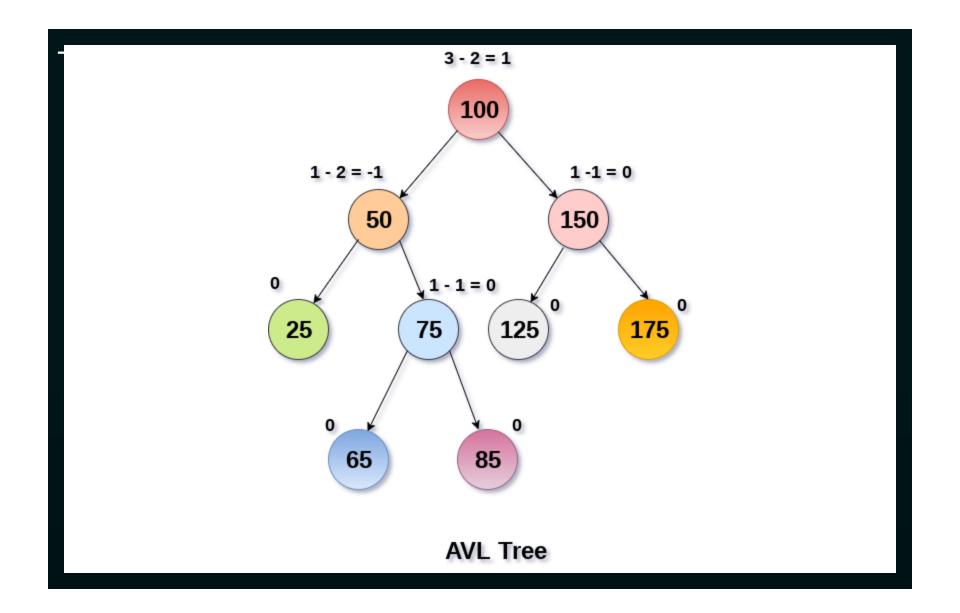
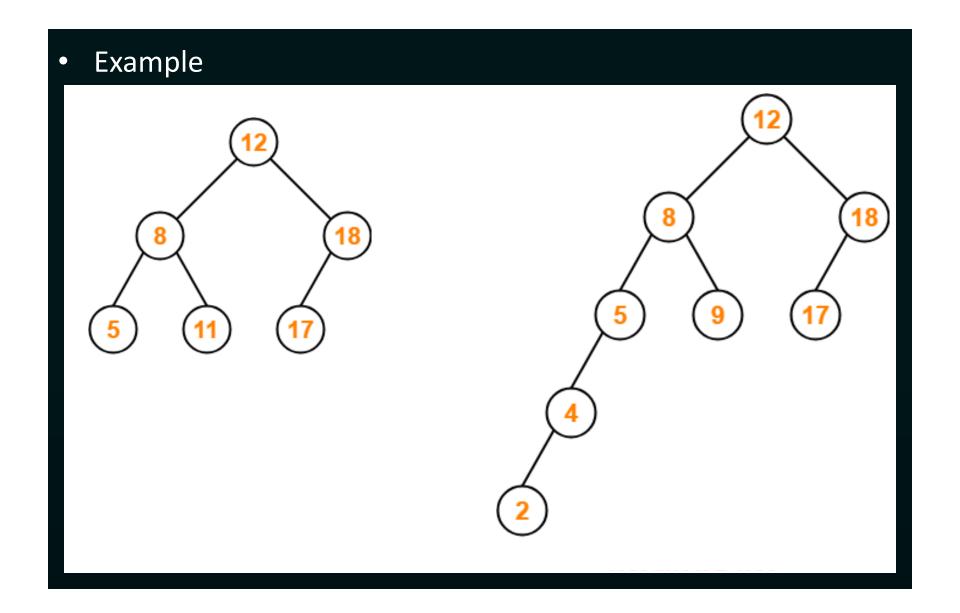
AVL Tree

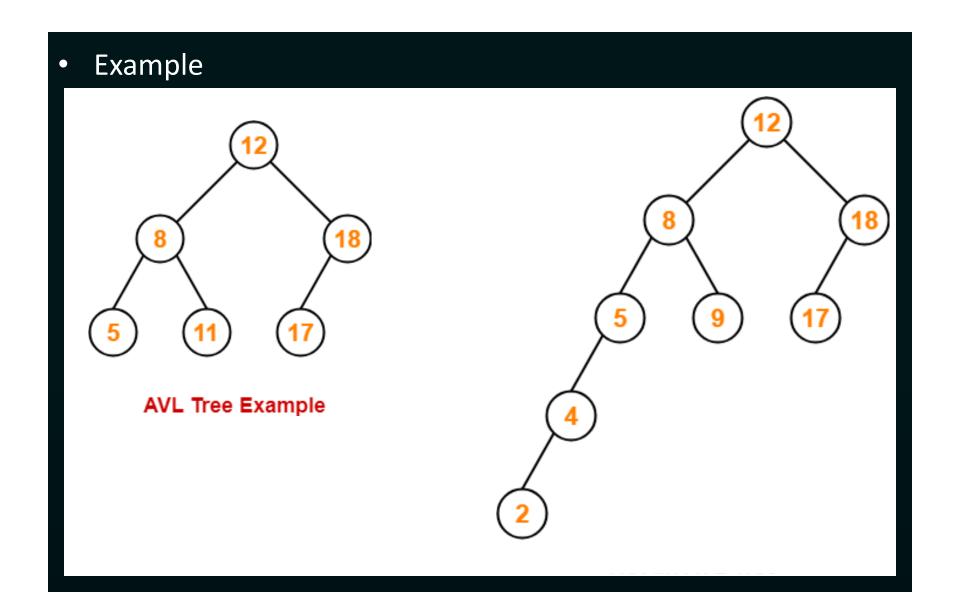
Introduction

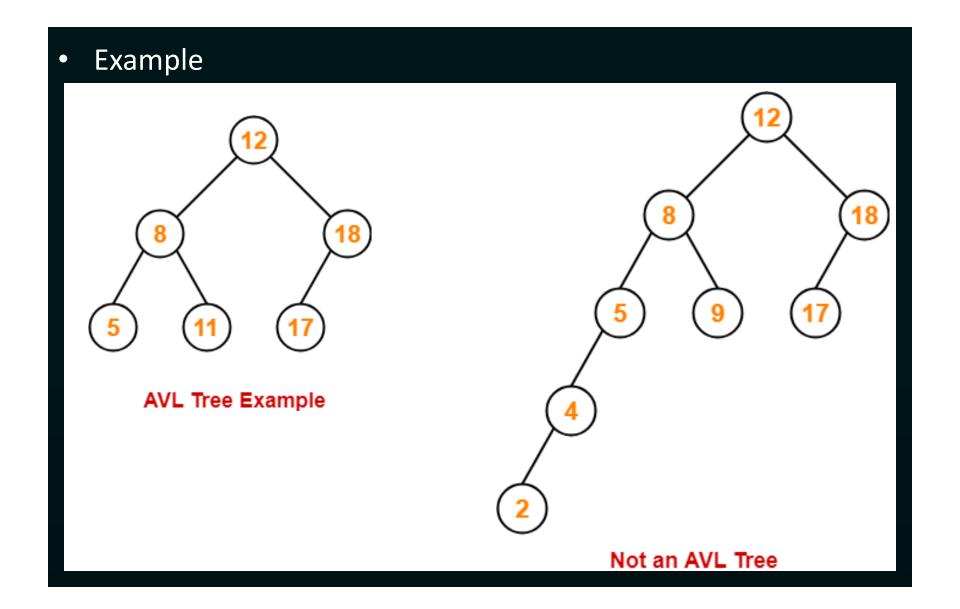
- AVL Tree
 - Special kind of binary search tree
 - Invented by Adelson, Velsky and Landis in 1962
 - Defined as height-balanced binary search tree
 - Each node is associated with a balance factor
 - Calculated by subtracting the height of its right sub-tree from that of its left sub-tree
 - Balance Factor (k) = height(left(k)) height(right(k))
 - Tree is said to be balanced
 - if balance factor of each node is either -1 or 0 or 1, otherwise, the tree will be unbalanced and need to be balanced

- In an AVL Tree
 - If balance factor of any node is -1, it means that the left sub-tree is one level lower than the right sub-tree
 - If balance factor of any node is 0, it means that the left sub-tree and right sub-tree contain equal height
 - If balance factor of any node is 1, it means that the left sub-tree is one level higher than the right sub-tree









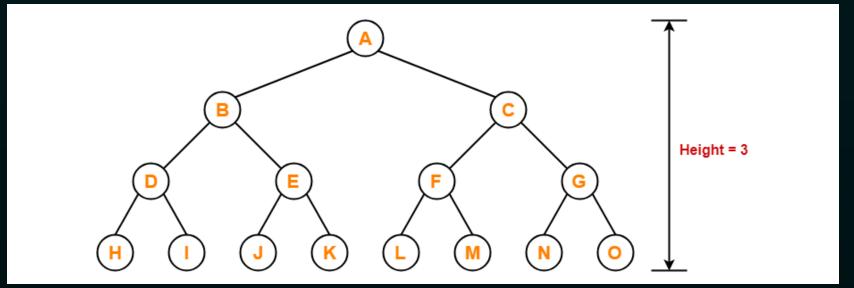
- Why AVL Trees?
 - AVL tree controls the height of the binary search tree by not letting it to be skewed
 - The time taken for all operations in a binary search tree of height h is O(h)
 - However, it can be extended to O(n) if the BST becomes skewed (i.e. worst case)
 - By limiting this height to log n, AVL tree imposes an upper bound on each operation to be O(log n) where n is the number of nodes

- AVL Tree Properties
 - Property-01: Max. no. possible number of nodes in AVL tree of height $H = 2^{H+1} 1$

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 - Example- Maximum possible number of nodes in AVL tree of height-3

- AVL Tree Properties
 - Property-01: Max. no. number of nodes in AVL tree of height $H = 2^{H+1} 1$
 - Example- Max. possible number of nodes in AVL tree of height 3

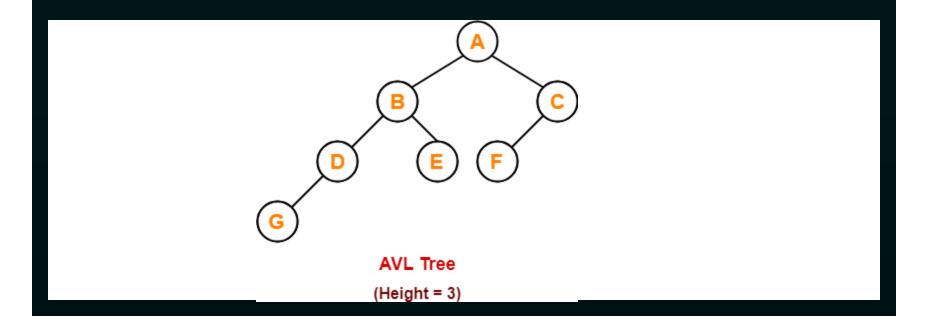
$$= 2^{3+1} - 1 => 16 - 1 => 15$$



- AVL Tree Properties
 - Property-02: Min. no. of nodes in AVL Tree of height H is given by a recursive relation

$$N(H) = N(H-1) + N(H-2) + 1$$

Given that N(0)=1, N(1)=2



- AVL Tree Properties
 - Property-03: Min possible height of AVL Tree using N nodes = [log₂N]

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 - Example: Min. possible height of AVL Tree using 8 nodes

- AVL Tree Properties
 - Property-03: Min possible height of AVL Tree using N nodes = [log₂N]
 - Example: Min. possible height of AVL Tree using 8 nodes

```
= [\log_2 8]
= [\log_2 2^3]
= [3\log_2 2]
= [3]
```

= 3

- Property-04: If there are N nodes in AVL Tree, its max. height can not exceed 1.44 log₂N
 - Worst case height of AVL Tree with N nodes = 1.44 log₂N
 - i.e. Worst case height of AVL Tree with n nodes = 1.44log₂n.

AVL Tree Operations

- Operations Search, Insertion and Deletion
- After performing any operation on AVL tree, the balance factor of each node is checked. There might be two cases possible

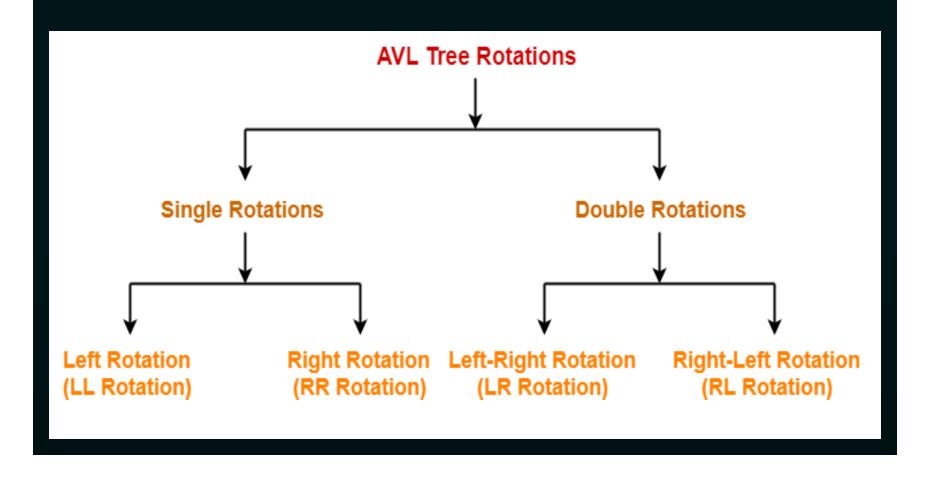
Case-01:

- After the operation, the balance factor of each node is either 0 or 1 or -1
- In this case, the AVL tree is considered to be balanced
- The operation is concluded.

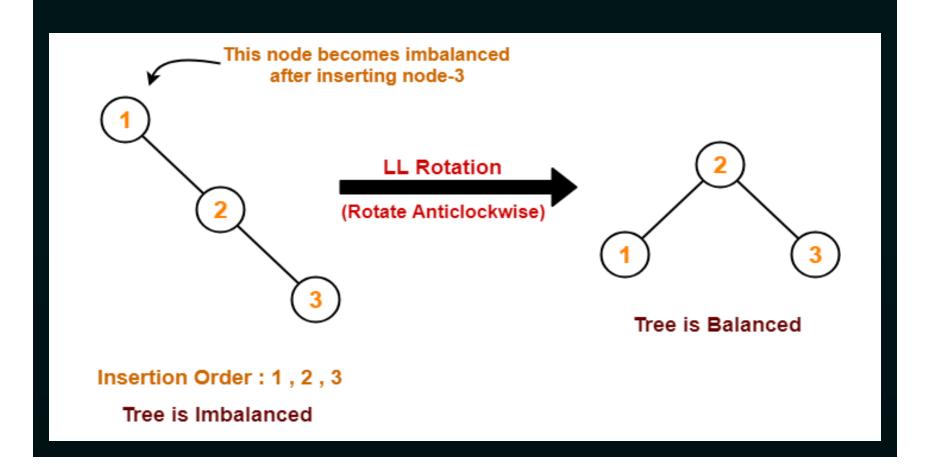
Case-02:

- After the operation, the balance factor of at least one node is not 0 or 1 or -1
- In this case, the AVL tree is considered to be imbalanced.
- Rotations are then performed to balance the tree.

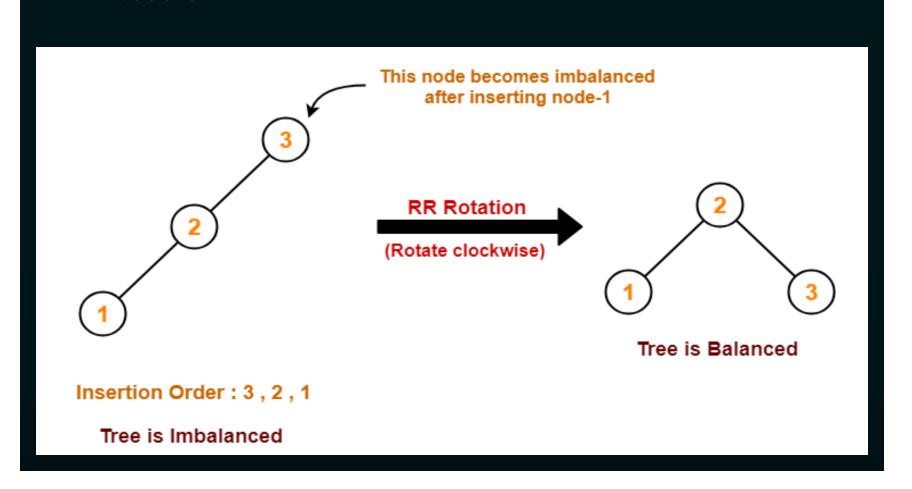
- AVL Tree Rotations
 - 4 kinds of rotations possible in AVL Trees



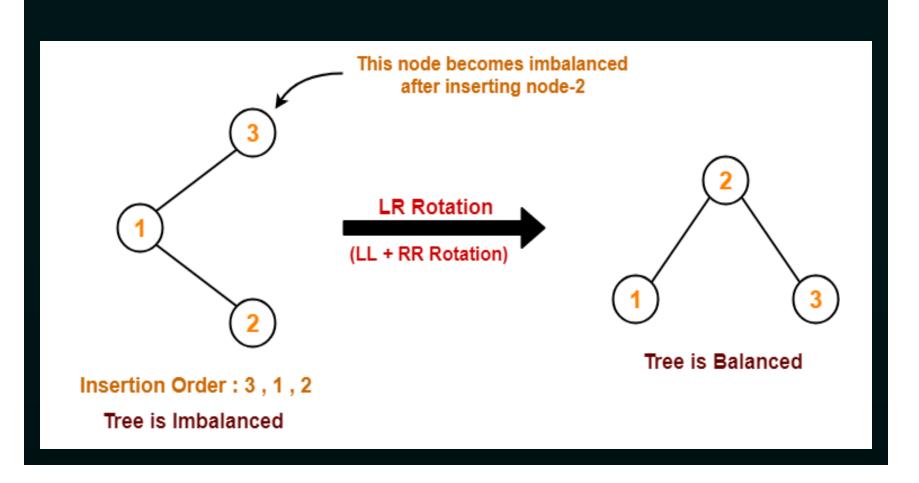
- Cases Of Imbalance
 - Case-01:



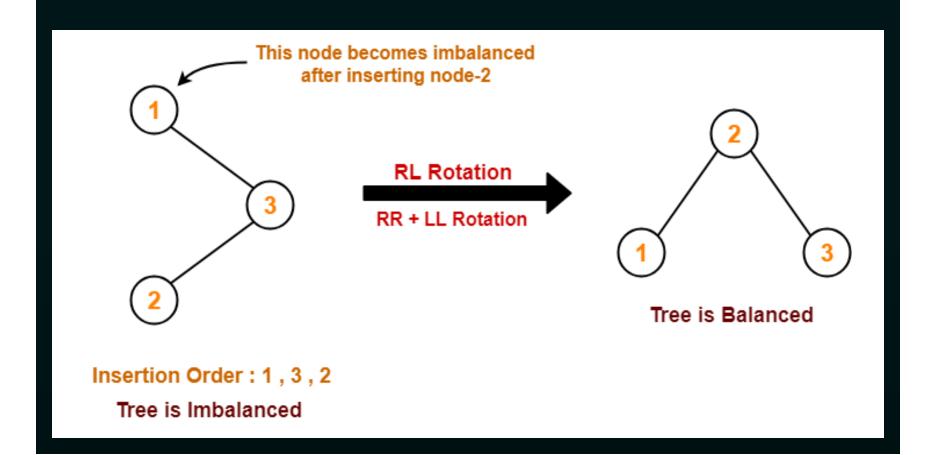
- Cases Of Imbalance
 - Case-02:



- Cases Of Imbalance
 - Case-03:



- Cases Of Imbalance
 - Case-04:



Insertion in AVL Tree

- To insert an element in the AVL tree, follow the following steps
 - Step 1: Insert the element in the AVL tree in the same way the insertion is performed in BST
 - Step 2: After insertion, check the balance factor of each node of the resulting tree
- Note-01:
 - Balance factor of only those nodes will be affected that lies on the path from the newly inserted node to the root node
- Note-02:
 - There is no need to check the balance factor of every node
 - Check the balance factor of only those nodes that lies on the path from the newly inserted node to the root node

Insertion in AVL Tree

• Note-03:

- After inserting an element in the AVL tree,
 - If tree becomes imbalanced, then there exists one particular node in the tree by balancing which the entire tree becomes balanced automatically.
 - To re-balance the tree, balance that particular node
- To find that particular node,
 - Check the balance factor of each node that is encountered while traversing the path from newly inserted node to the root node
 - The first encountered imbalanced node will be the node that needs to be balanced
- To balance that node,
 - Count three nodes in the direction of leaf node and use the concept of AVL tree rotations to re balance the tree

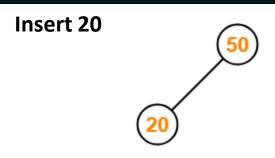
 Construct AVL Tree for the following sequence of numbers-50, 20, 60, 10, 8, 15, 32, 46, 11, 48

Insert 50

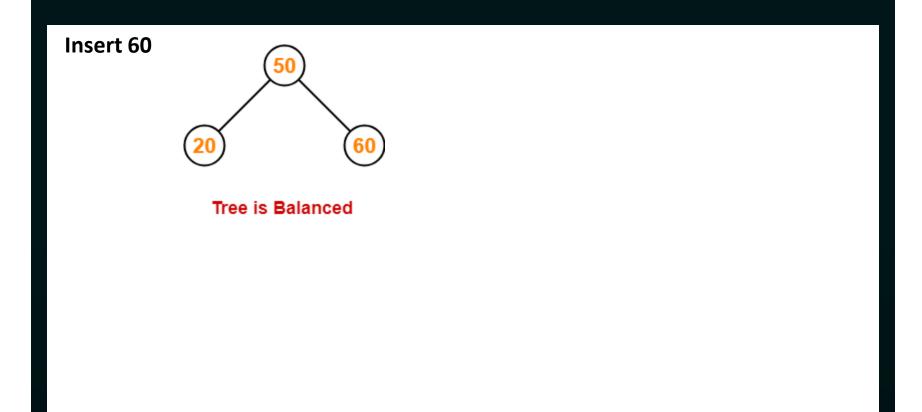


Tree is Balanced

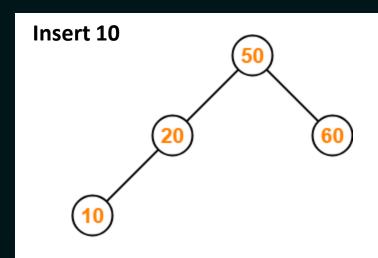
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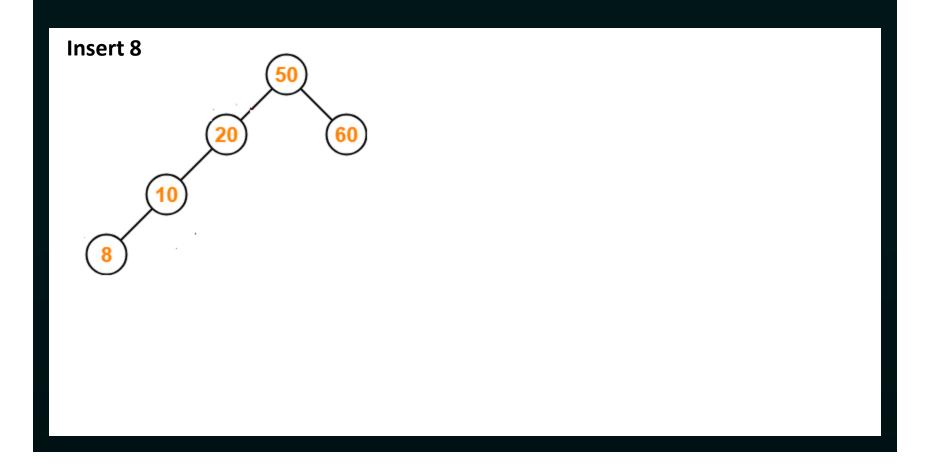
Tree is Balanced



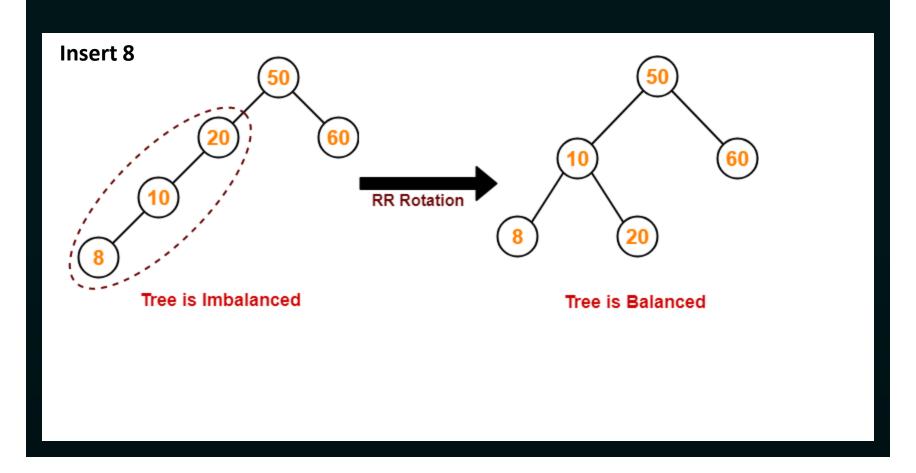
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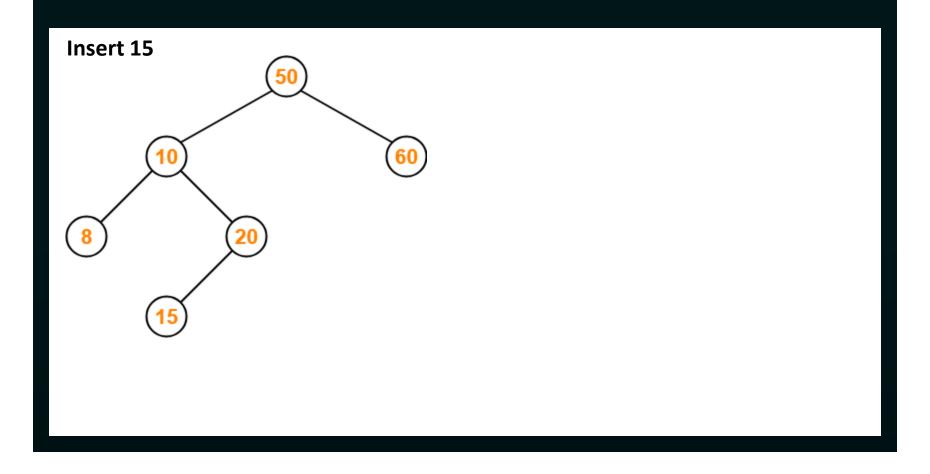


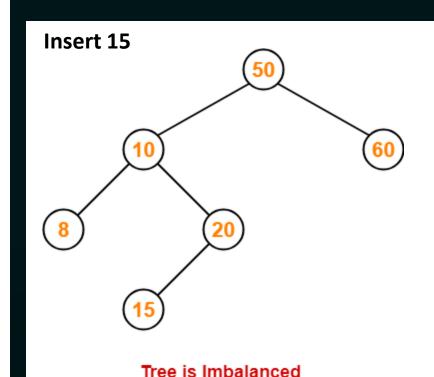
Tree is Balanced

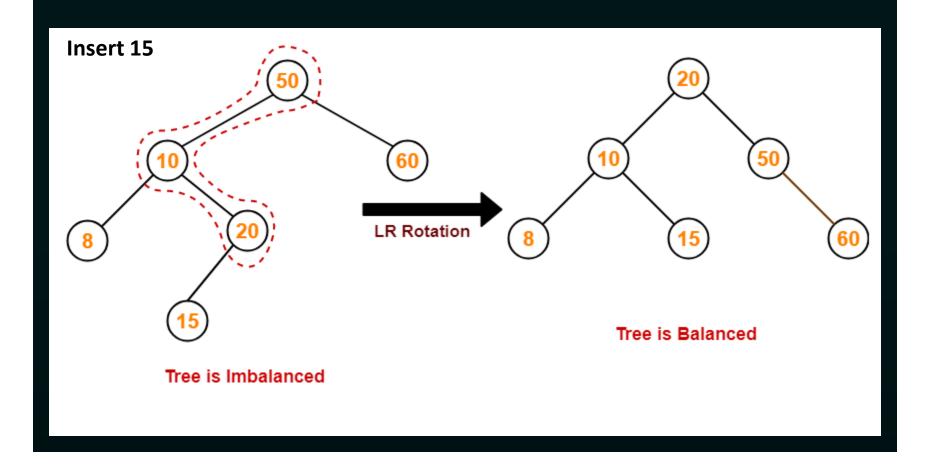


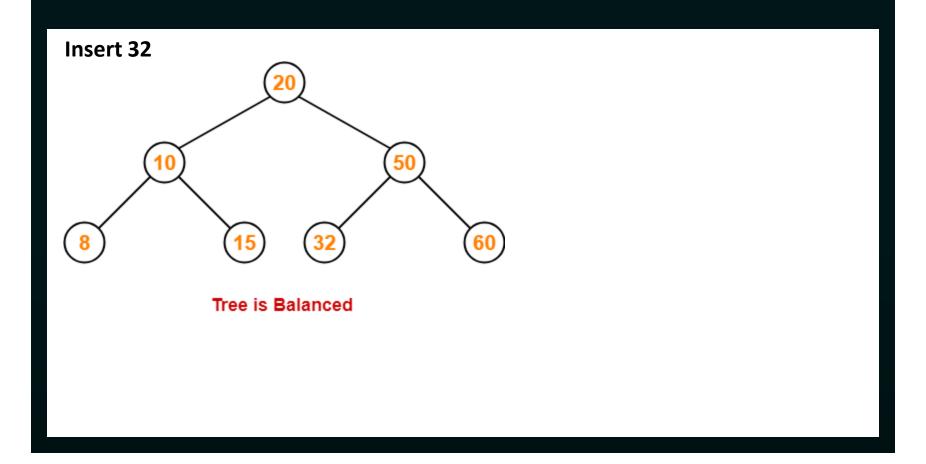


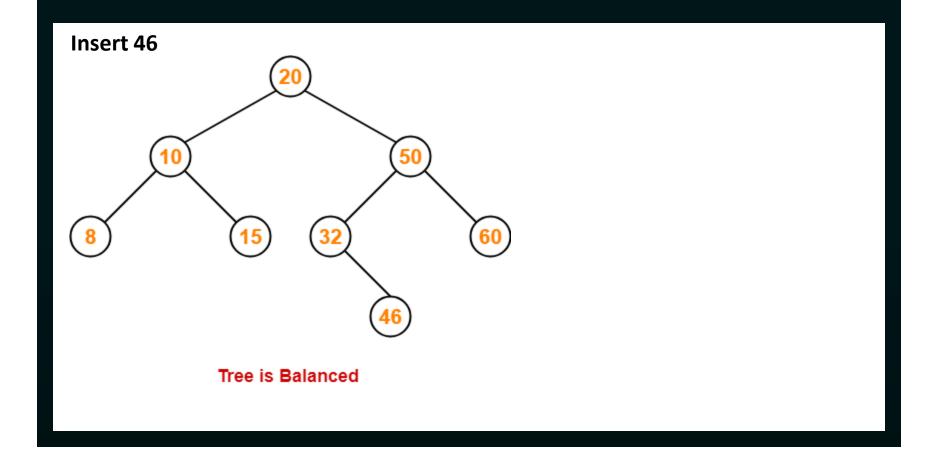


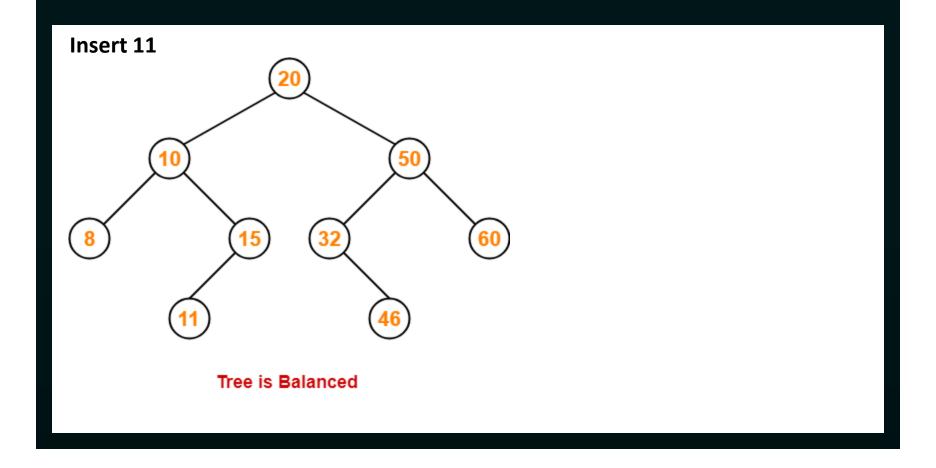


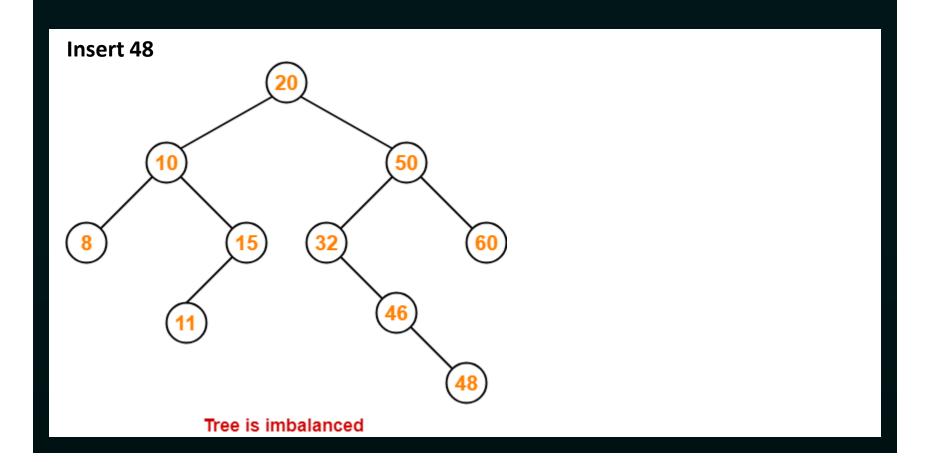


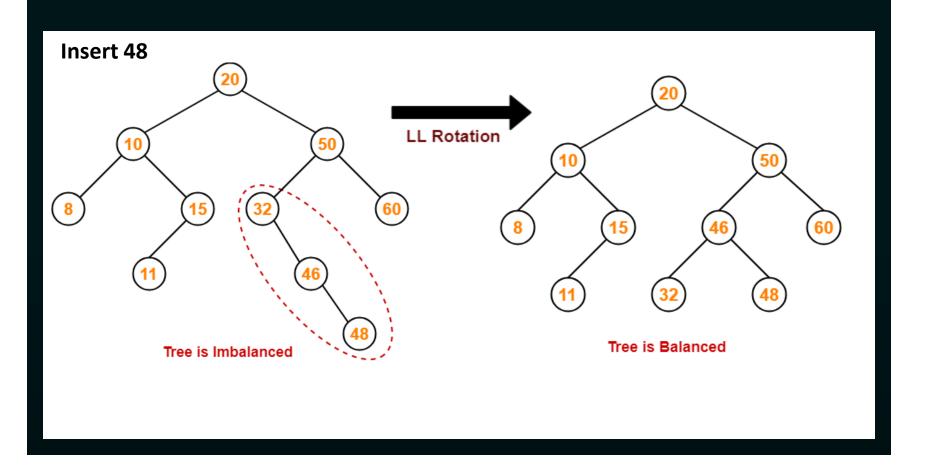


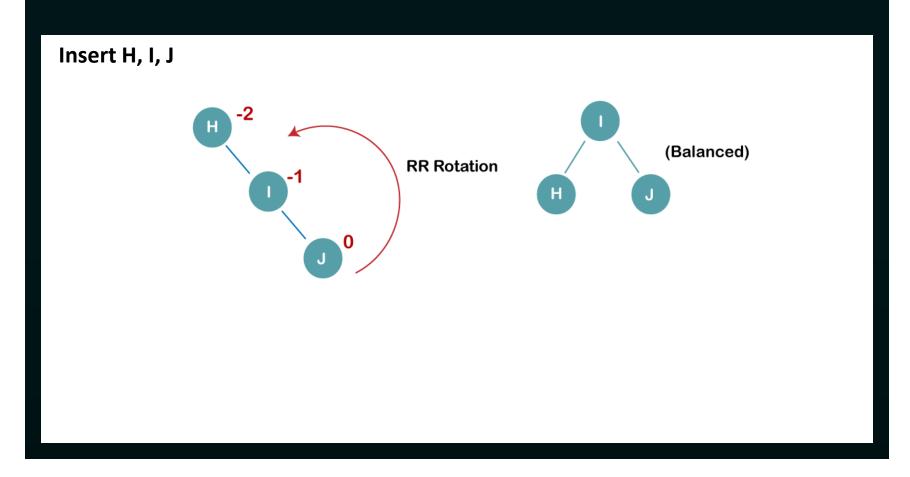












Construct an AVL tree having the following elements
 H, I, J, B, A, E, C, F, D, G, K, L

Insert B, A

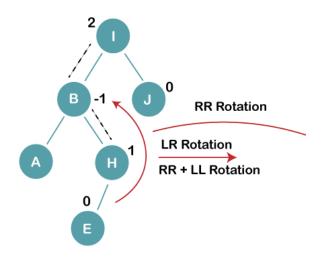
2
H

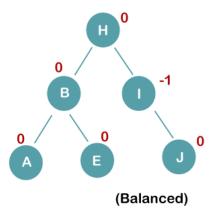
2
H

(Balanced)

Construct an AVL tree having the following elements
 H, I, J, B, A, E, C, F, D, G, K, L

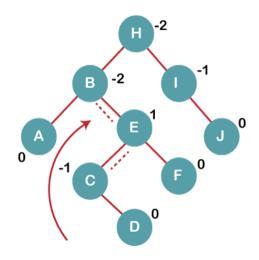
Insert E

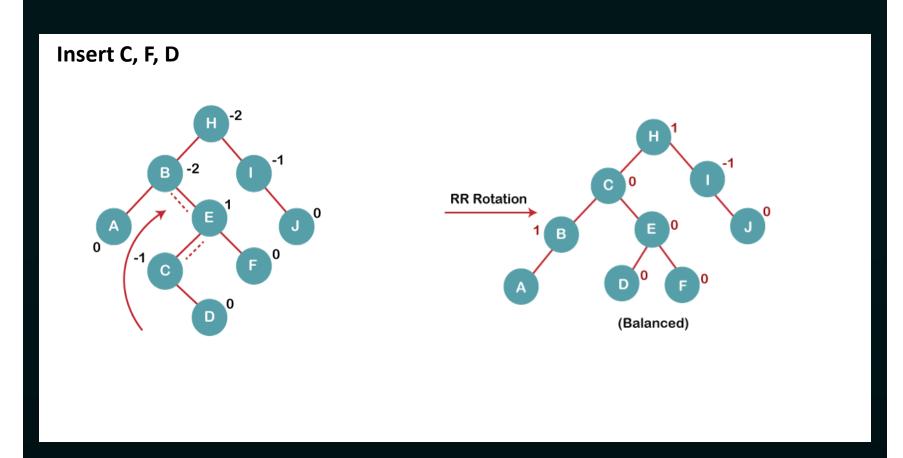




Construct an AVL tree having the following elements
 H, I, J, B, A, E, C, F, D, G, K, L

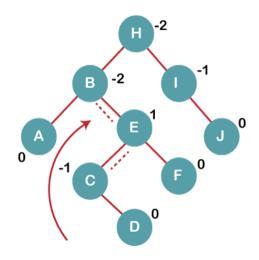
Insert C, F, D

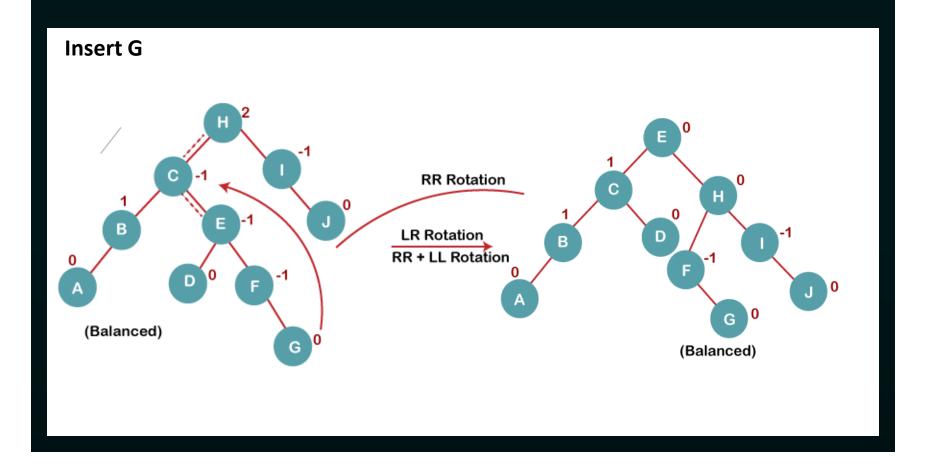


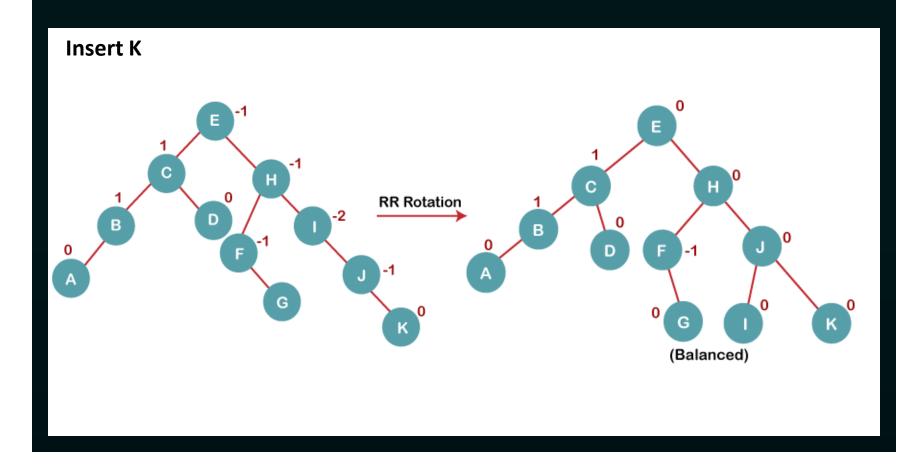


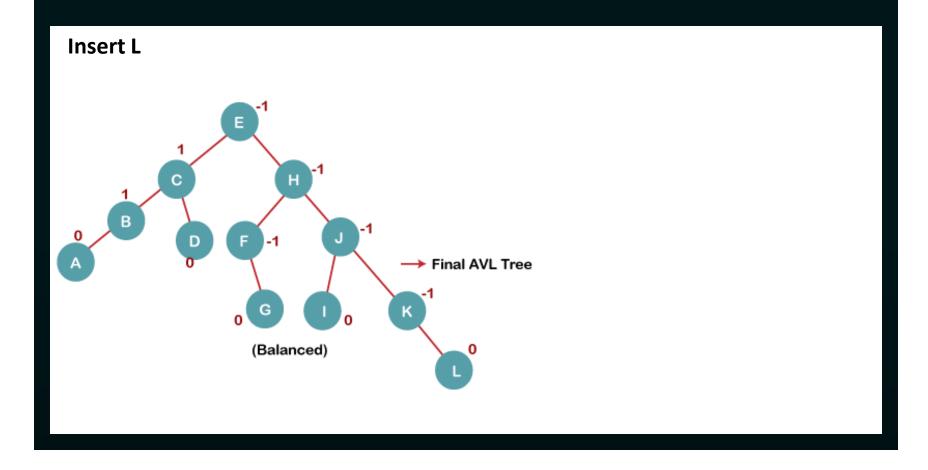
Construct an AVL tree having the following elements
 H, I, J, B, A, E, C, F, D, G, K, L

Insert C, F, D





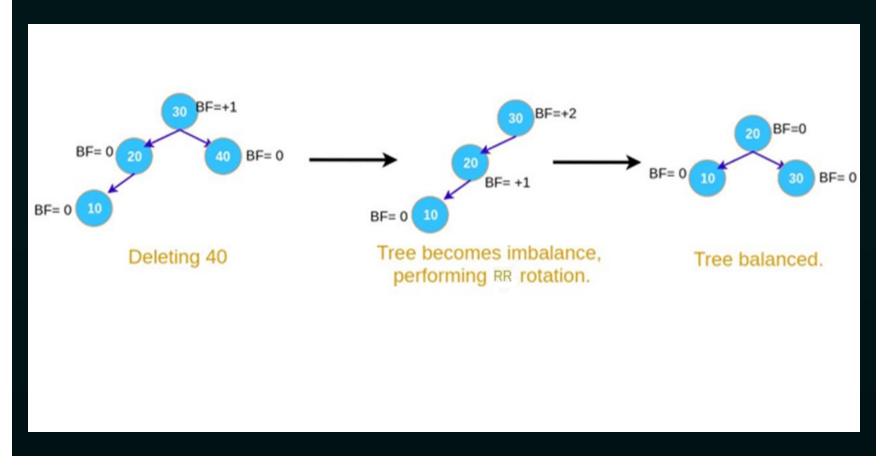




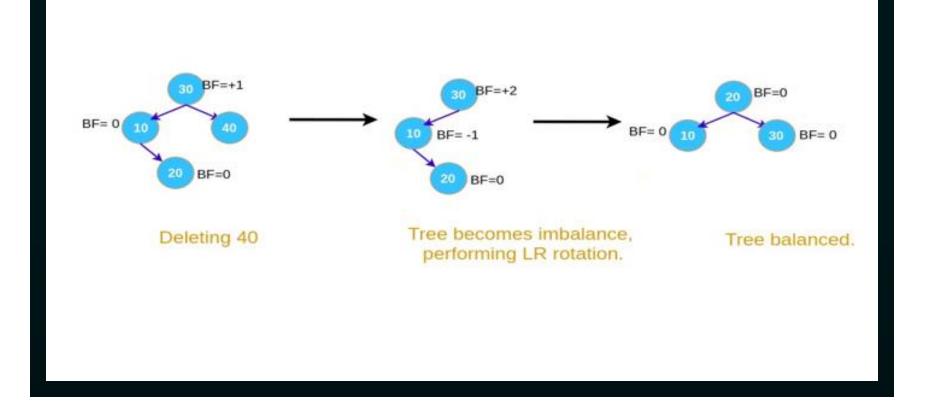
Deletion in AVL Tree

- Deletion in AVL Trees
 - Done using the same logic as in binary search tree
 - However, after deletion, restructure the tree, if needed, to maintain its balanced height
- Steps of Deletion
 - Step 1: Find the element in the tree.
 - Step 2: Delete the node, as per the BST Deletion.
 - Step 3: Two cases are possible:-
 - Case 1: Deleting from the right subtree.
 - Case 2: Deleting from left subtree.

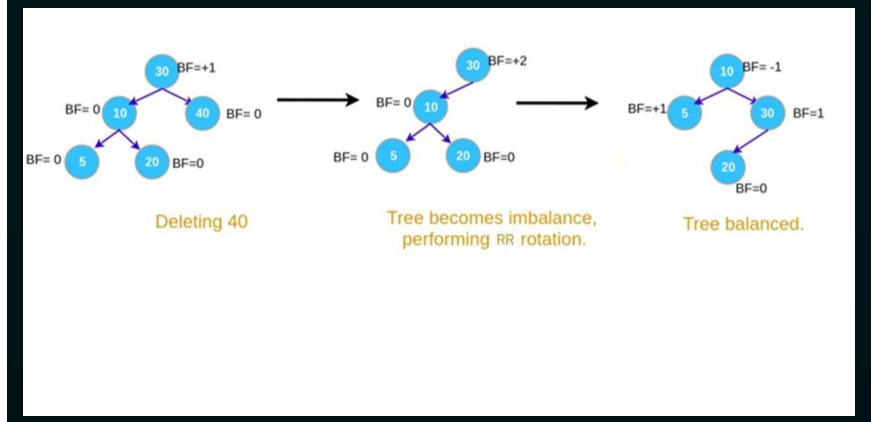
- Case 1: Deleting from the right subtree
 - 1A. If BF(node)=+2 and BF(node->left-child) = +1, perform RR rotation



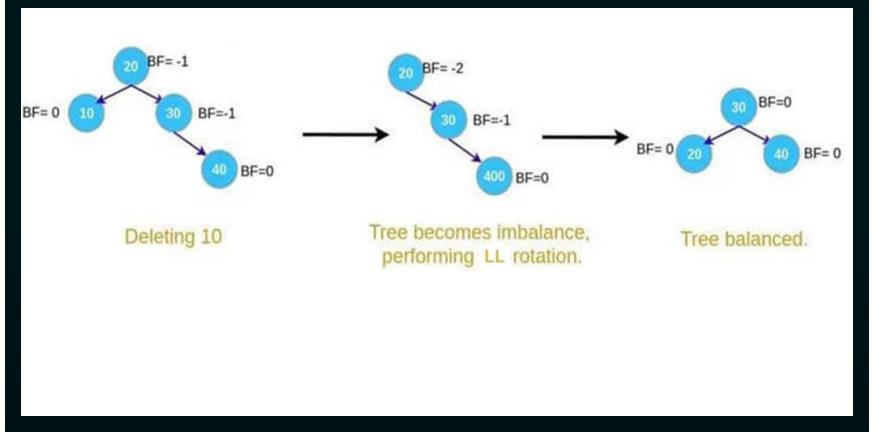
- Case 1: Deleting from the right subtree
 - 1B. If BF(node)=+2 and BF(node->left-child)=-1, perform LR rotation.



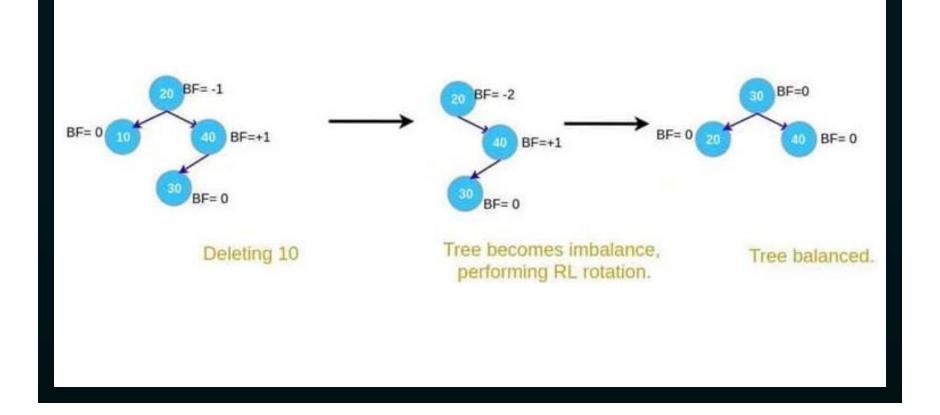
- Case 1: Deleting from the right subtree
 - 1C. If BF(node)=+2 and BF(node->left-child)= 0, perform RR rotation.



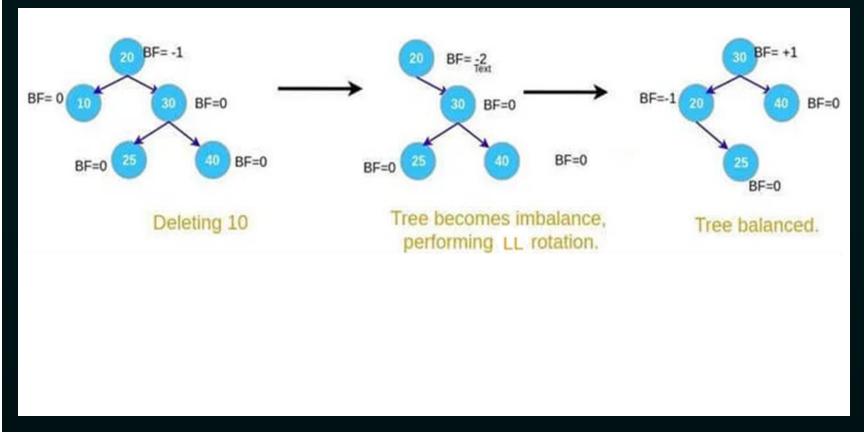
- Case 2: Deleting from the left subtree
 - 2A. If BF(node)=-2 and BF(node->right-child) =-1, perform LL rotation.



- Case 2: Deleting from the left subtree
 - 2B. If BF(node)=-2 and BF(node->right-child)=+1, perform RL rotation.



- Case 2: Deleting from the left subtree
 - 2C. If BF(node)=-2 and BF(node->right-child)=0, perform LL rotation.



- Advantages of AVL Trees
 - The height of the AVL tree is always balanced. The height never grows beyond log N, where N is the total number of nodes in the tree.
 - It gives better search time complexity when compared to simple Binary Search trees.
 - AVL trees have self-balancing capabilities.

- Summary
 - AVL trees are self-balancing binary search trees.
 - Balance factor is the fundamental attribute of AVL trees
 - The balance factor of a node is defined as the difference between the height of the left and right subtree of that node.
 - The valid values of the balance factor are -1, 0, and +1.
 - The insert and delete operation require rotations to be performed after violating the balance factor.
 - The time complexity of insert, delete, and search operation is O(log N).
 - AVL trees follow all properties of Binary Search Trees.
 - The left subtree has nodes that are lesser than the root node. The right subtree has nodes that are always greater than the root node.
 - AVL trees are used where search operation is more frequent compared to insert and delete operations.