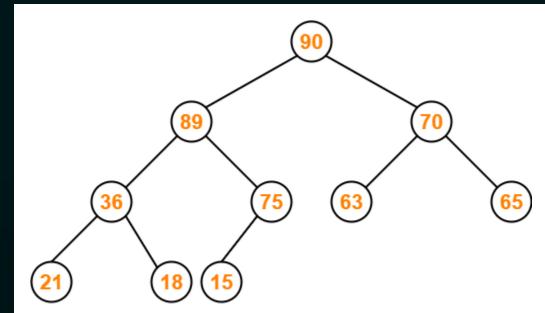
Introduction

- Heap
 - A specialized data structure and has special characteristics.
 - May be implemented using a n-ary tree.
- Binary Heap
 - Binary Tree with the following two properties
 - 1. Ordering Property: Elements in the heap are arranged in specific order.
 - Two types of heaps- min heap & max heap
 - 2. Structural Property: Binary heap is an almost complete binary tree
 - All its levels completely filled except possibly the last level
 - The last level is strictly filled from left to right

Types of Binary Heap Depending on the arrangement of elements, a binary heap may be of following two types-Types of Binary Heap Max Heap Min Heap

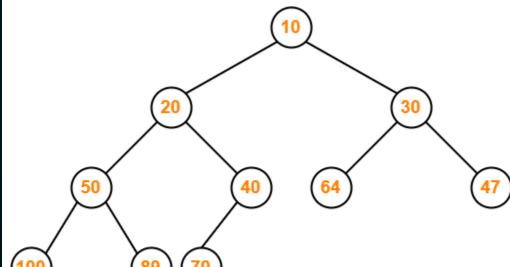
- Max Heap
 - In max heap, every node contains greater or equal value element than its child nodes
 - Thus, root node contains the largest value element



- Every node contains greater or equal value element than its child nodes.
- It is an almost complete binary tree with its last level strictly filled from left to right.

Max Heap Example

- Min Heap
 - In min heap, every node contains lesser value element than its child nodes
 - Thus, root node contains the smallest value element.

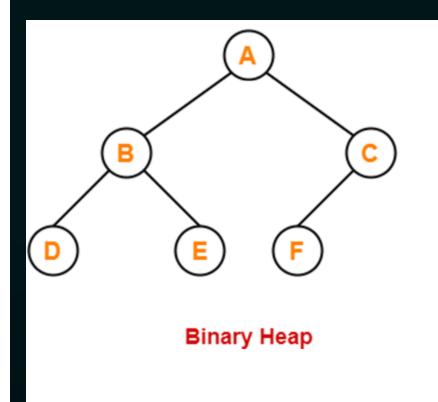


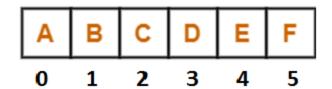
- Every node contains lesser value element than its child nodes.
- It is an almost complete binary tree with its last level strictly filled from left to right.

Min Heap Example

Array Representation of Binary Heap

A binary heap is typically represented as an Array





Array Representation

- For a node present at index 'i' of the array Arr[]
 - If indexing starts with 0,
 - Its parent node present at array location = Arr[(i-1)/2]
 - Its left child node present at array location = Arr[2i+1]
 - Its right child node present at array location = Arr[2i+2]
 - If indexing starts with 1,
 - Its parent node present at array location = Arr [[i/2]]
 - Its left child node present at array location = Arr [2i]
 - Its right child node present at array location = Arr[2i+1]

• Note-01:

- Level order traversal technique may be used to achieve the array representation of a heap tree
- Array representation of a heap never contains any empty indices in between

Note-02:

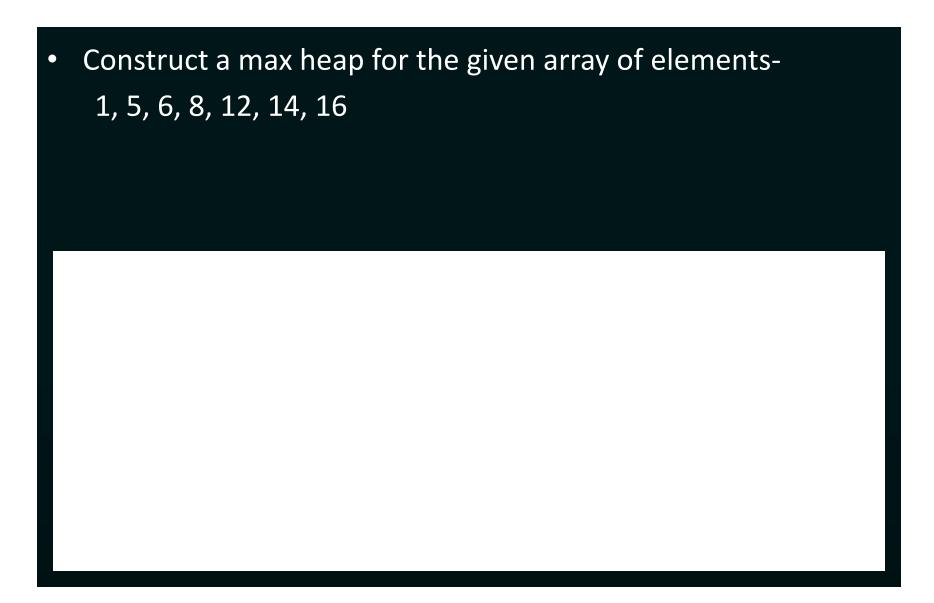
- Given an array representation of a binary heap,
 - If all the elements are in descending order, then heap is definitely a max heap
 - If all the elements are in ascending order, then heap is definitely a min heap

- Consider a binary max-heap implemented using an array.
 Which one of the following array represents a binary max-heap?
 - a) 25, 14, 16, 13, 10, 8, 12
 - b) 25, 12, 16, 13, 10, 8, 14
 - c) 25, 14, 12, 13, 10, 8, 16
 - d) 25, 14, 13, 16, 10, 8, 12

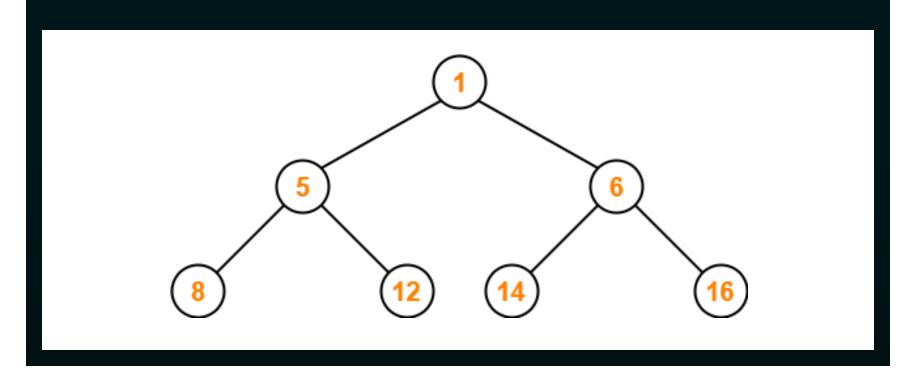
Max Heap Construction

- Given an array of elements, the steps involved in constructing a max heap
 - Step-01: Convert the given array of elements into an almost complete binary tree
 - Step-02: Ensure that the tree is a max heap
 - Check that every non-leaf node contains a greater or equal value element than its child nodes
 - If there exists any node that does not satisfies the ordering property of max heap, swap the elements
 - Start checking from a non-leaf node with the highest index (bottom to top and right to left)

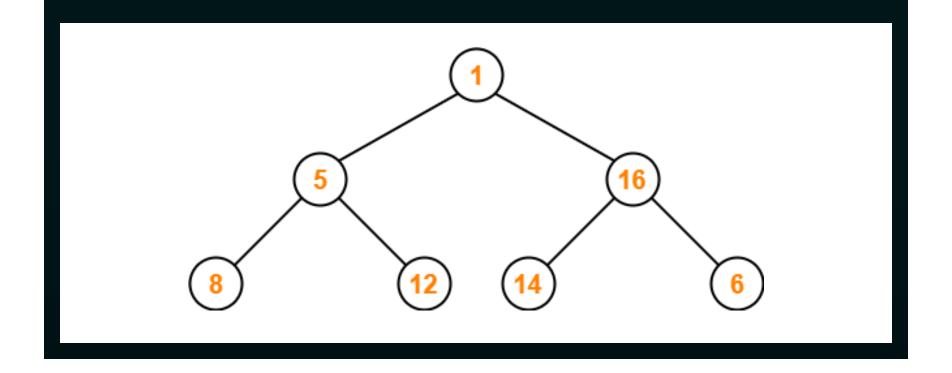
Example: Max Heap Construction



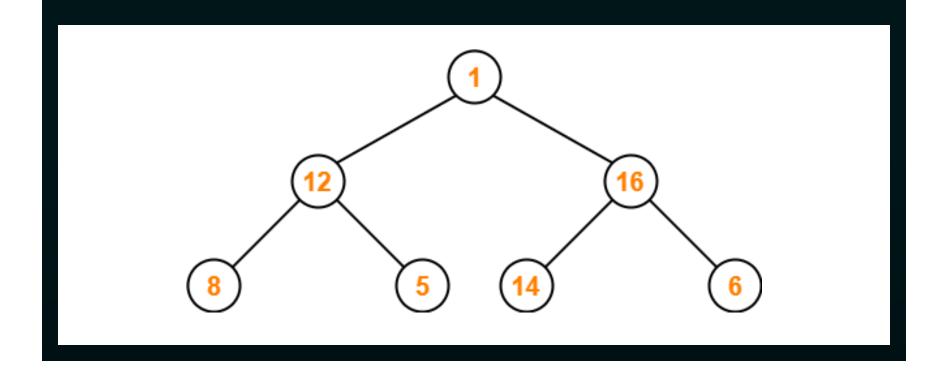
- Construct a max heap for the given array of elements 1, 5, 6, 8, 12, 14, 16
- Step-01: Convert the given array of elements into an almost complete binary tree



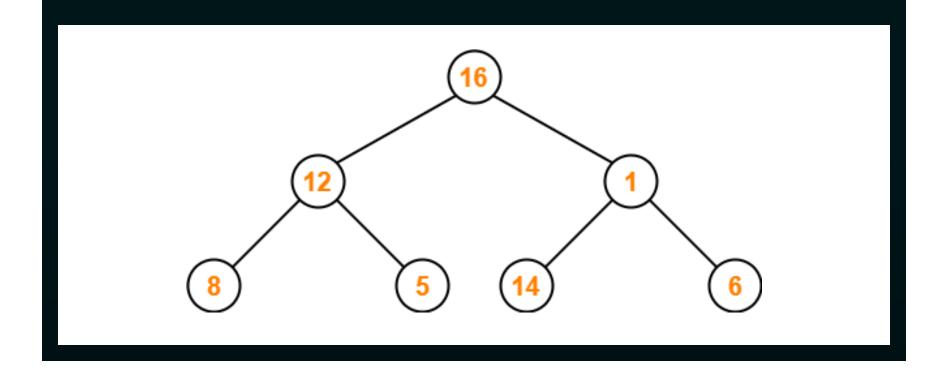
- Construct a max heap for the given array of elements 1, 5, 6, 8, 12, 14, 16
- Step-02:



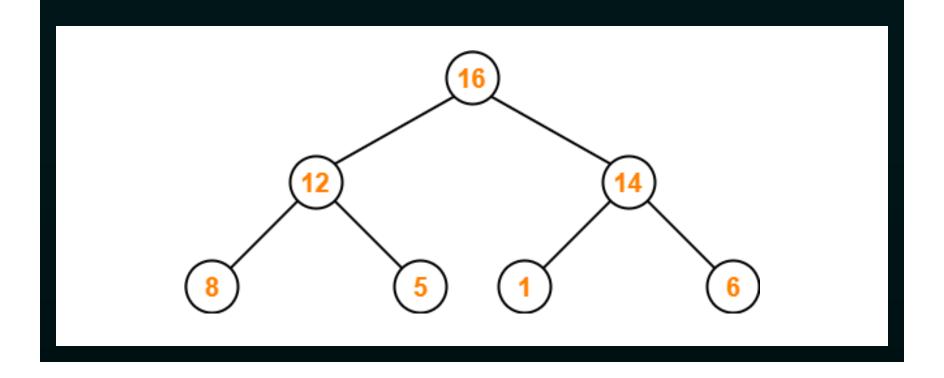
- Construct a max heap for the given array of elements 1, 5, 6, 8, 12, 14, 16
- Step-03:



- Construct a max heap for the given array of elements 1, 5, 6, 8, 12, 14, 16
- Step-04:



- Construct a max heap for the given array of elements 1, 5, 6, 8, 12, 14, 16
- Step-04:



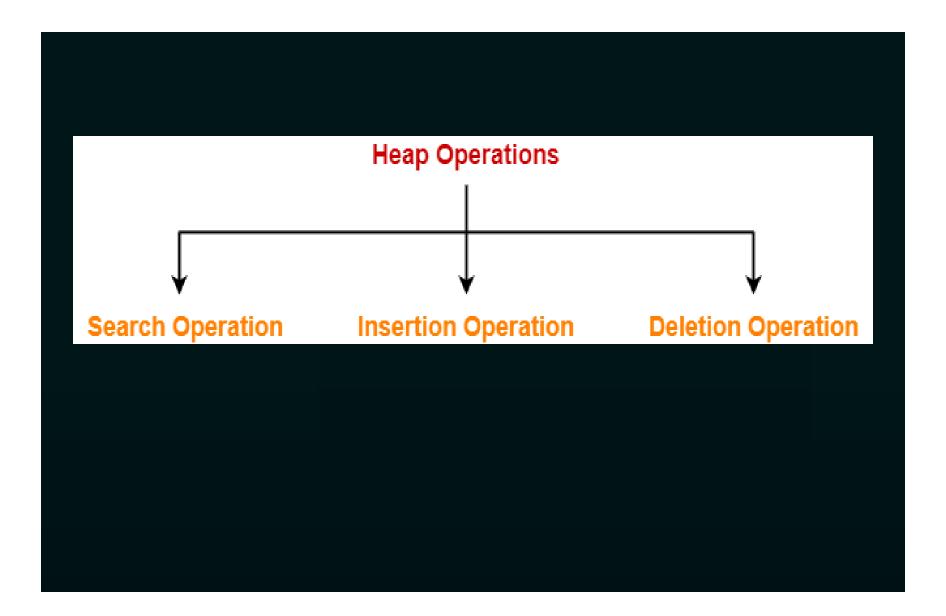
```
   Algorithm for Building Heap

  BUILD_HEAP(A, n)
      last_parent = floor(n/2)-1;
      for( i = last_parent; i>=0; i--)
              MAX_HEAPIFY(A, i);
```

Algorithm for Maintaining the heap property

```
MAX HEAPIFY(A, i)
 largest = i;
 I = 2*i +1;
 r = 2*i + 2;
 if (I < A.heapsize and A[I] > A[i])
         largest = I;
 if (r < A.heapsize and A[r] > A[largest])
         largest = r;
  if (largest != i)
         exchange A[i] with A[largest]
         MAX_HEAPIFY(A, largest)
```

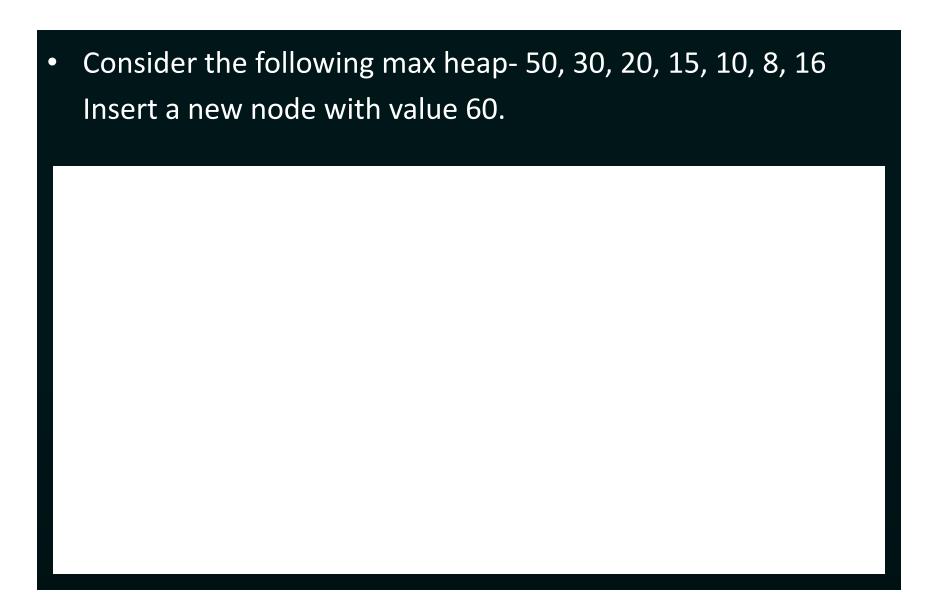
Heap Operations

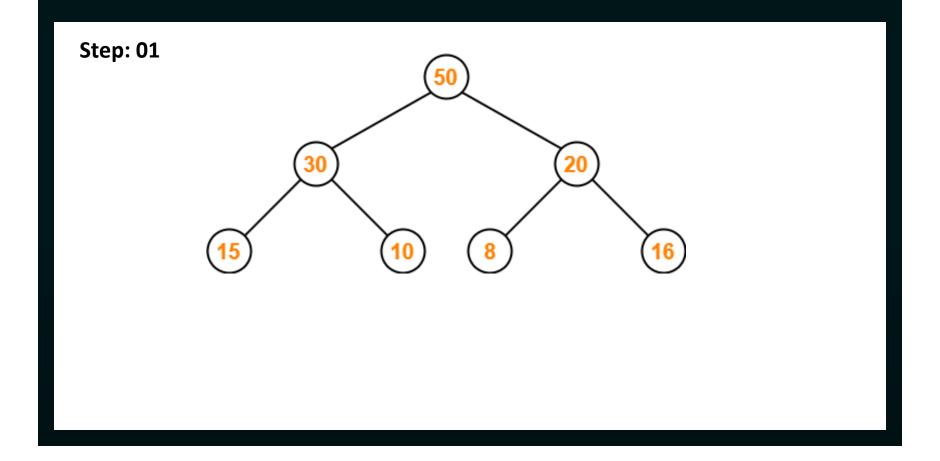


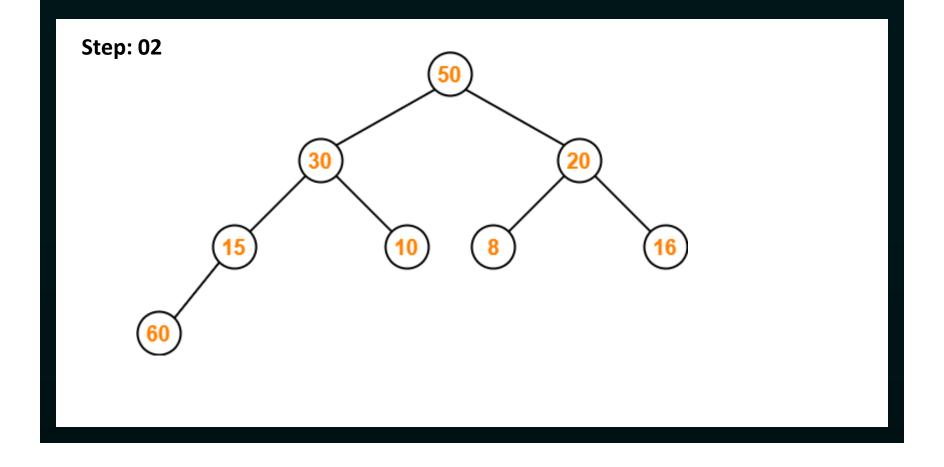
Insertion Operation(heapify_up)

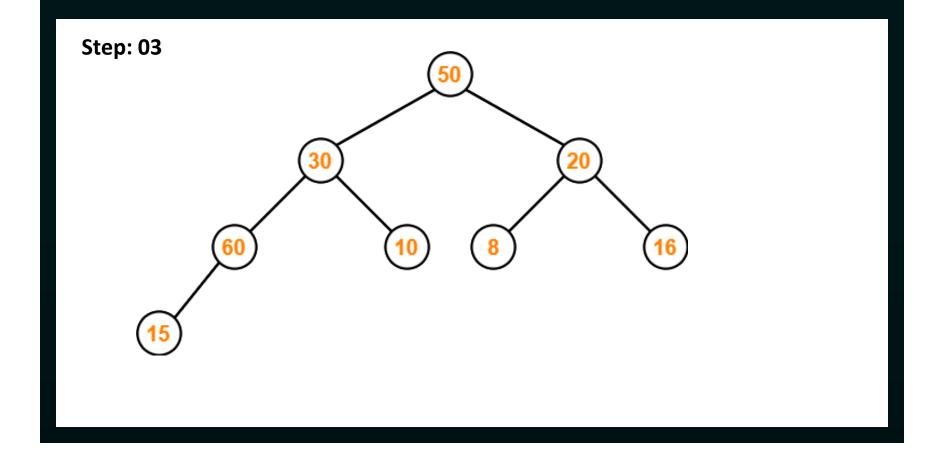
- Step-01:
 - Insert the new element as a next leaf node from left to right
- Step-02:
 - Ensure that the tree remains a max heap
 - Check that every non-leaf node contains a greater or equal value element than its child nodes
 - If there exists any node that does not satisfies the ordering property of max heap, swap the elements
 - Start checking from a non-leaf node with the highest index (bottom to top and right to left)

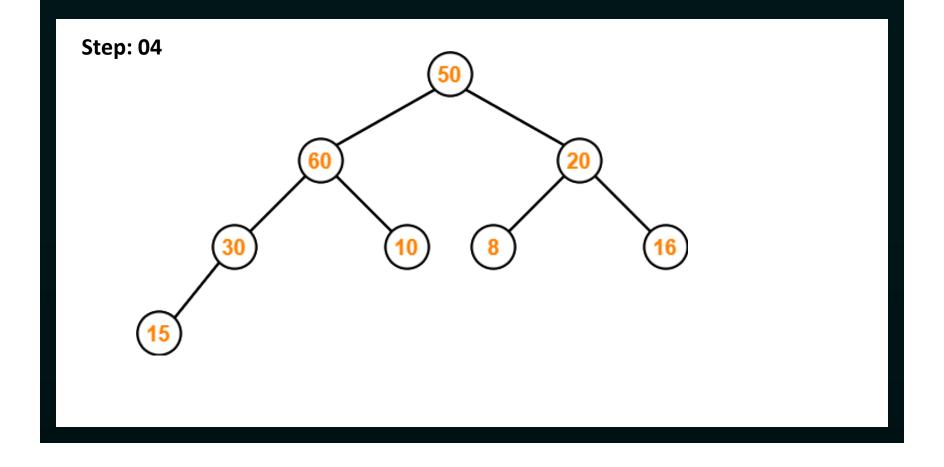
Example: Insertion of Node

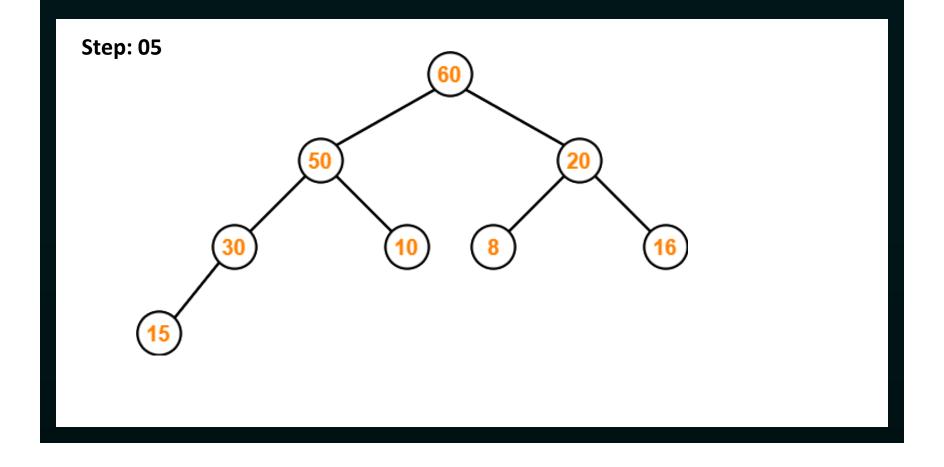








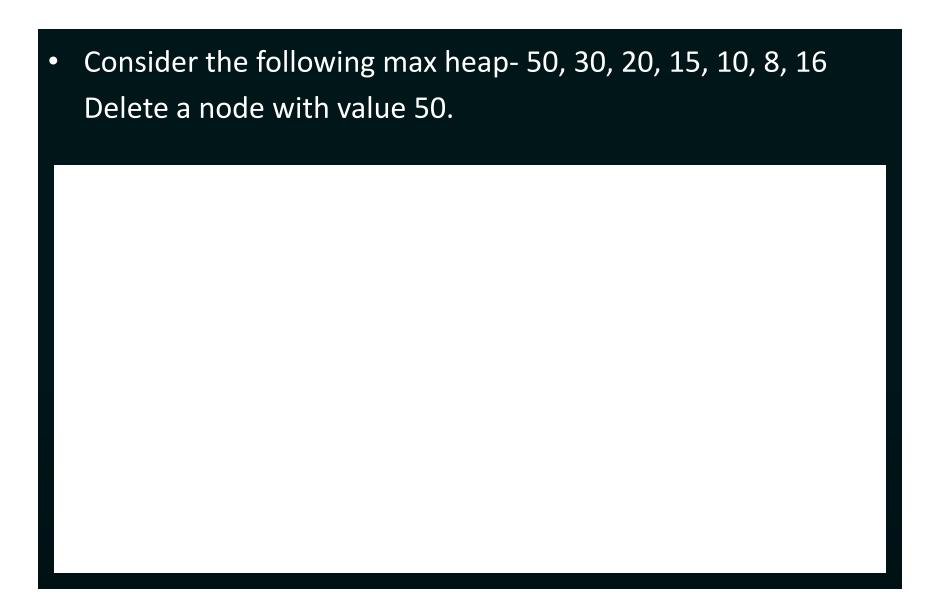


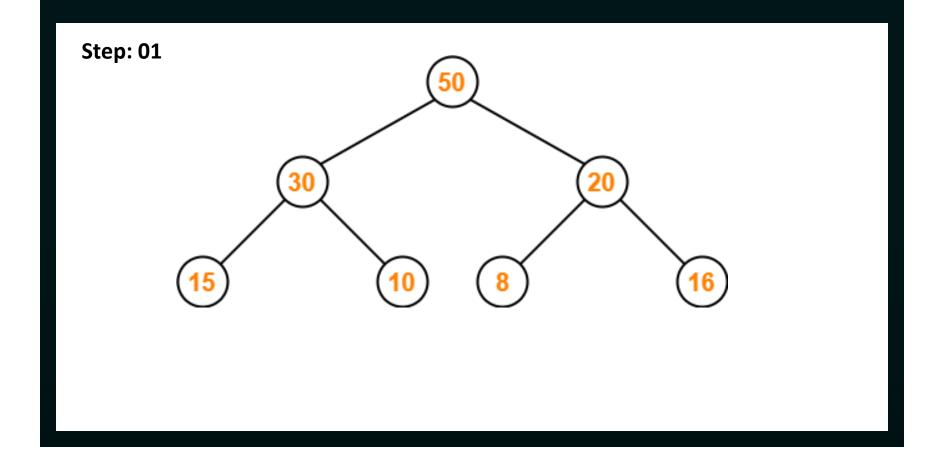


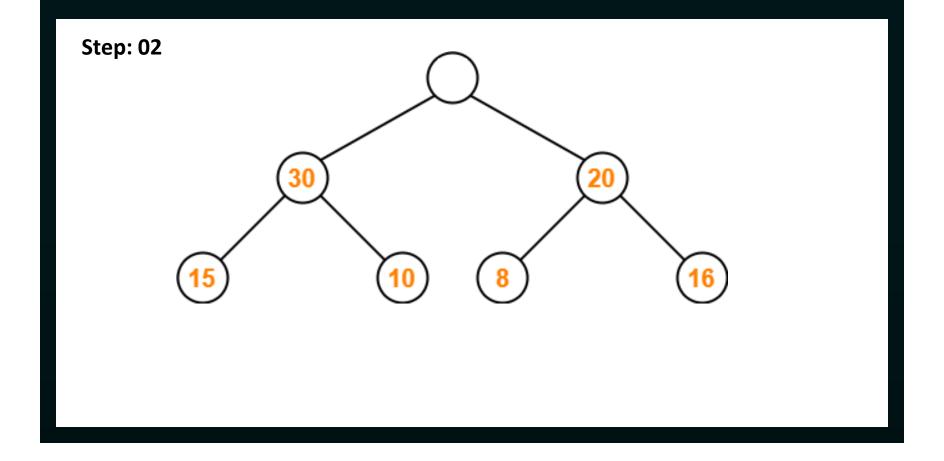
Deletion Operation

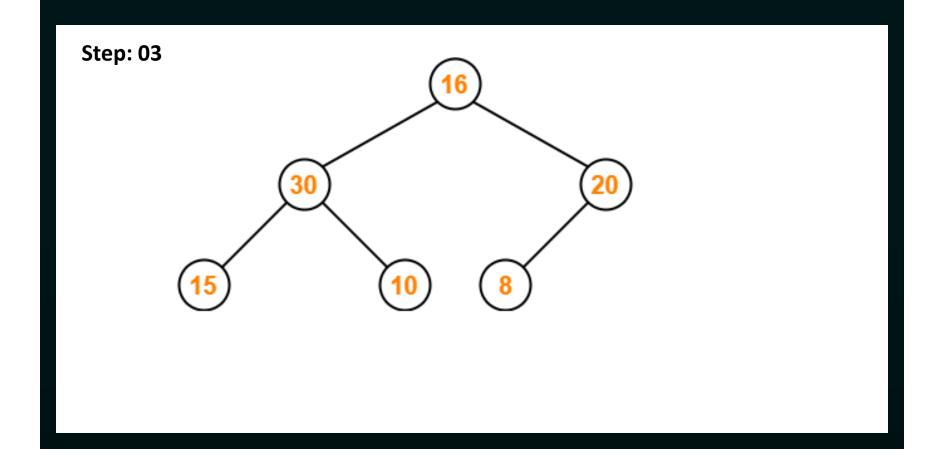
- Case-01: Deletion of Last Node
 - Just remove/disconnect the last leaf node from the heap
- Case-02: Deletion of Some Other Node disturbs heap properties
 - Step-01:
 - Delete the desired element from the heap tree.
 - Pluck last node and put in place of the deleted node
 - Step-02:
 - Ensure that the tree remains a max heap

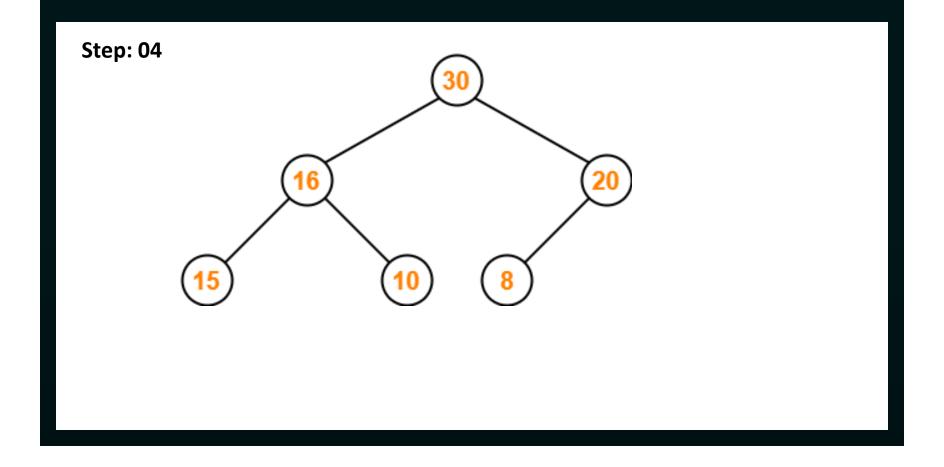
Example: Deletion of a Node





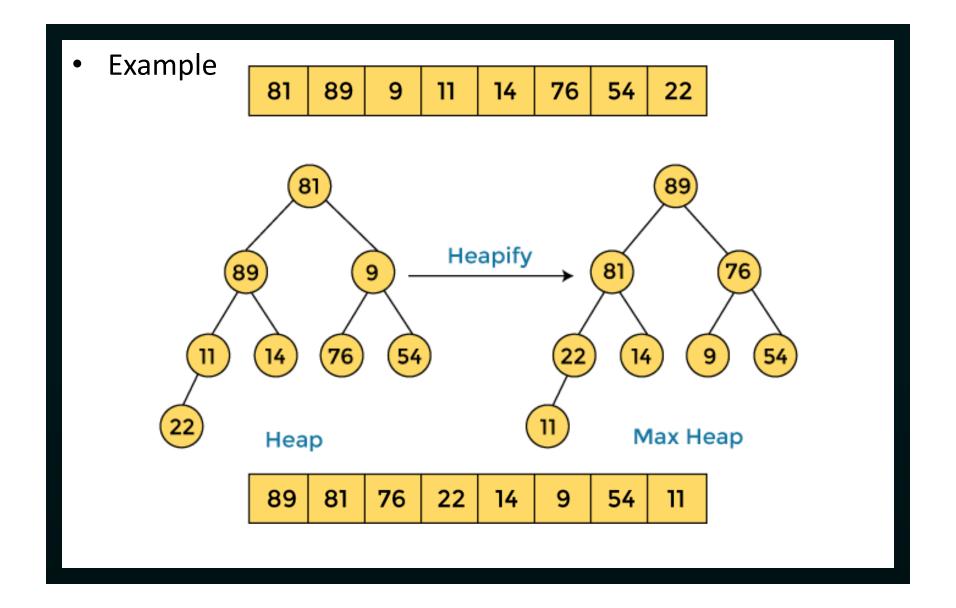


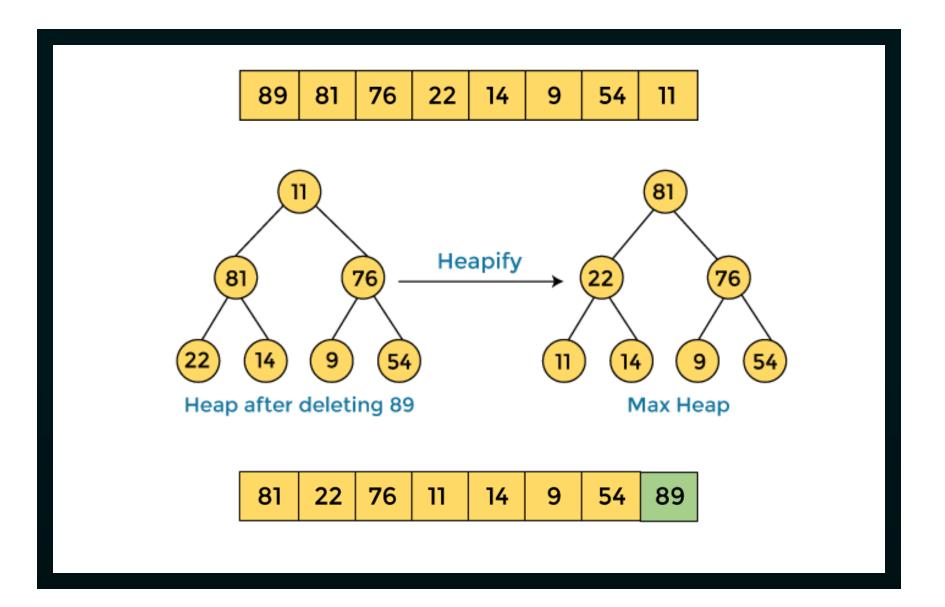


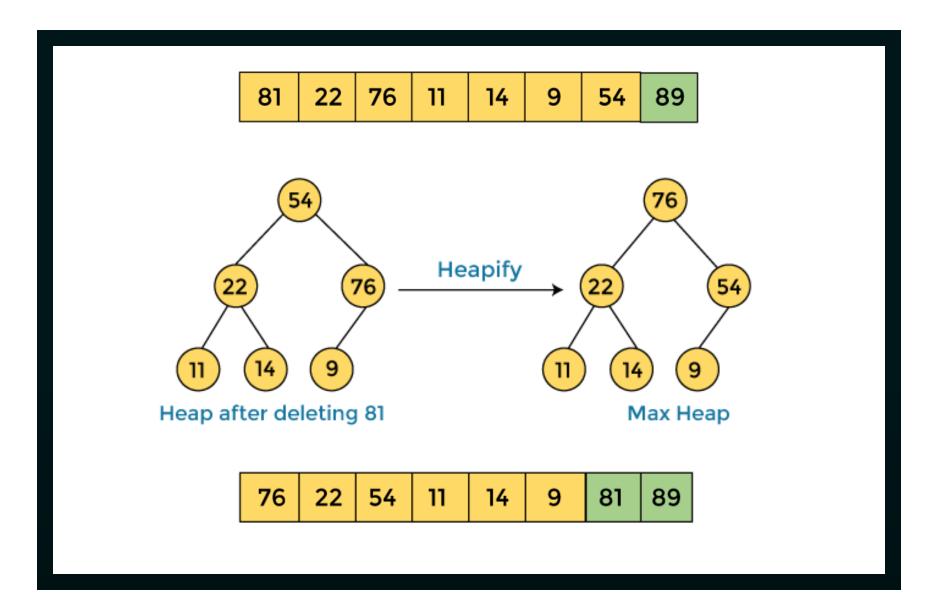


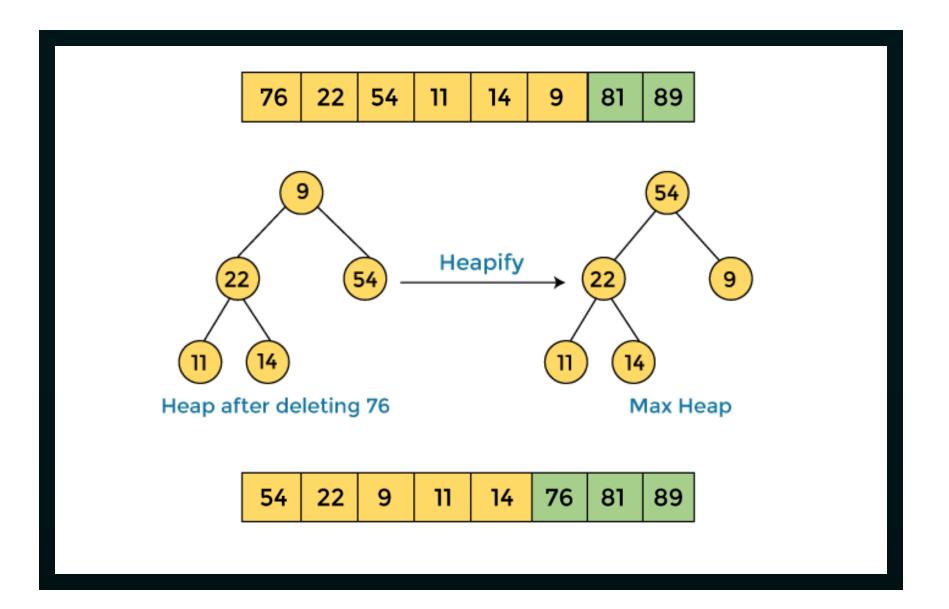
Heap Sort

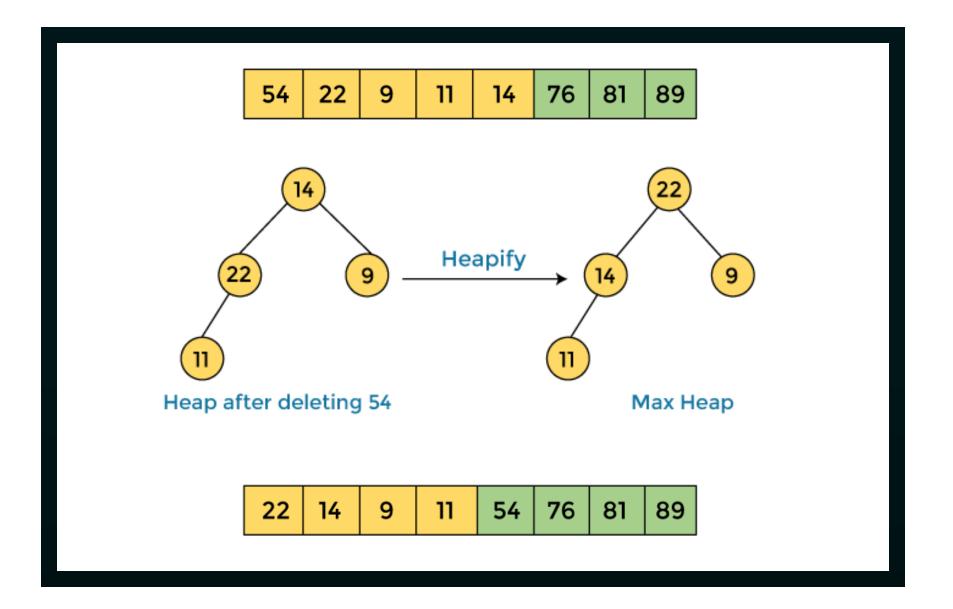
- Heap sort processes the elements by creating the min-heap or max-heap using the elements of the given array
- Concept of heap sort
 - Eliminate the elements one by one from the heap part of the list
 - and then insert them into the sorted part of the list
- Heap sort recursively performs two main operations
 - Build a heap H, using the elements of array
 - Repeatedly delete the root element of the heap formed in 1st step

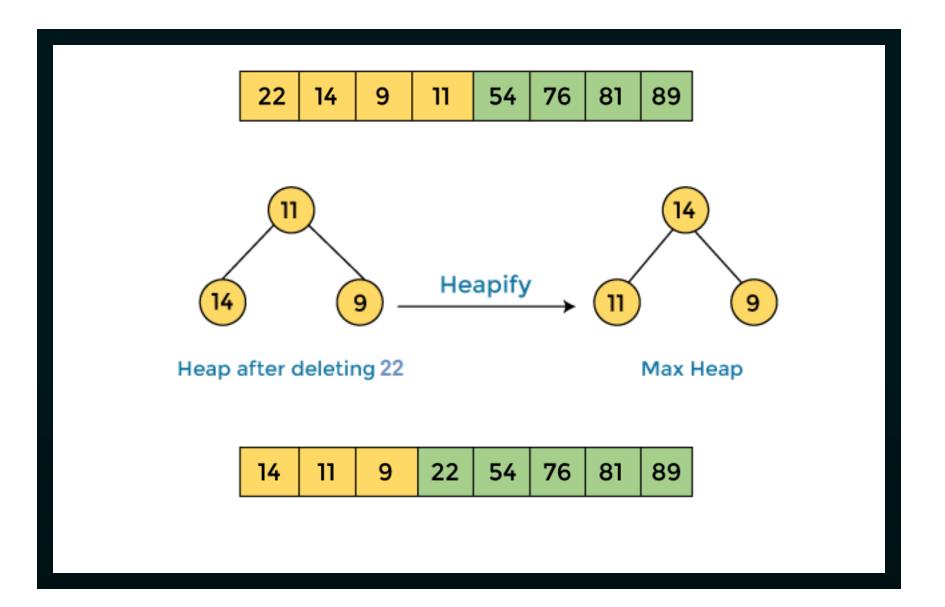


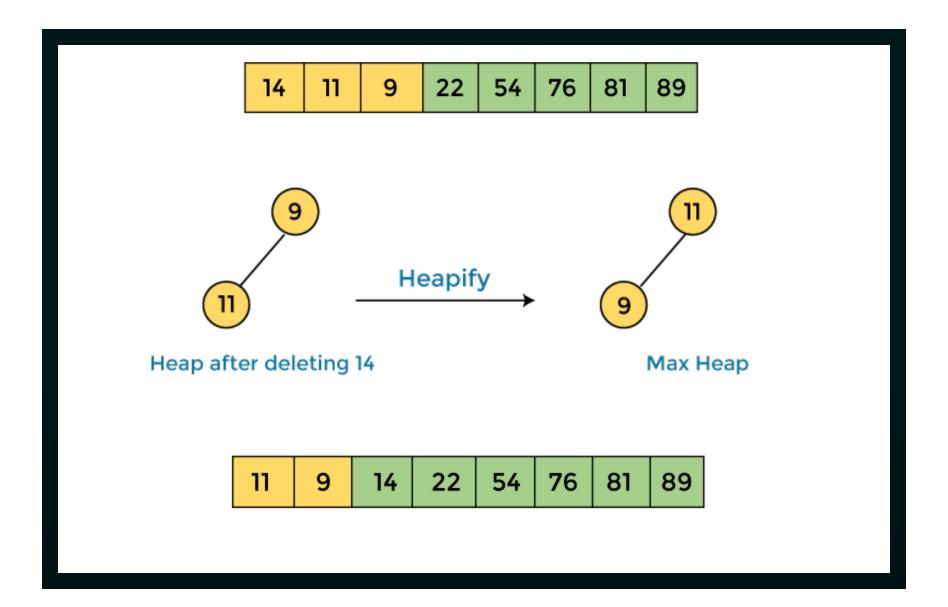


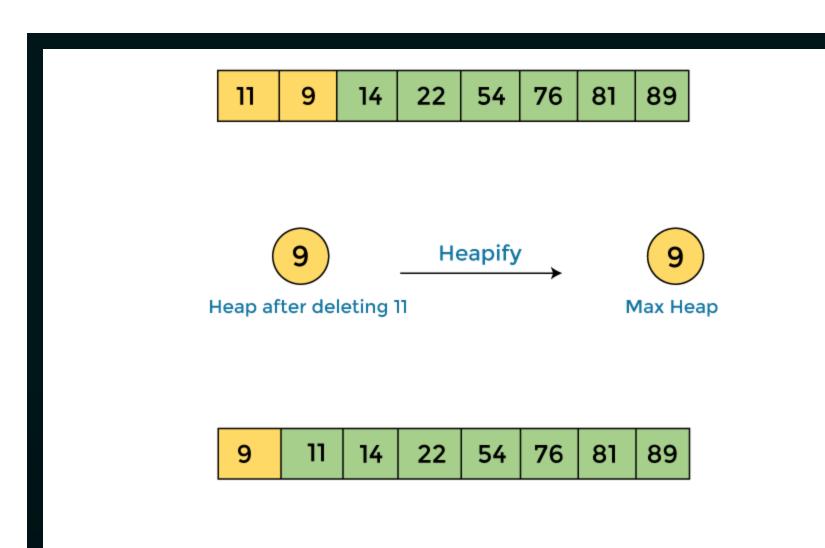


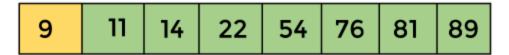














9 | 11 | 14 | 22 | 54 | 76 | 81 | 89

```
void heapSort(int a[], int n)
for (int i = n/2-1; i>=0; i--)
  heapify(a, n, i);
  // One by one extract an element from heap
for (int i = n-1; i>=0; i--)
    /* Move current root element to end*/
   // swap a[0] with a[i]
    int temp = a[0];
    a[0] = a[i];
    a[i] = temp;
    heapify(a, i, 0);
```

- Time Complexity
 - Best Case Complexity O(n log n)
 - Average Case Complexity O(n log n)
 - Worst Case Complexity O(n log n)