

Queues and Singly Linked Lists

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Queue

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Queue ADT

- **Queues** store arbitrary objects.
- **Insertions:** At the **end** of the queue.
- **Removals:** From the **front** of the queue.
- The queue has a **head** and a **tail**.



Queue ADT (Cont.)

Main Operations:

- **ENQUEUE**(Q, x): Inserts an element at the end.
- **DEQUEUE**(Q): Removes and returns the element at the front.



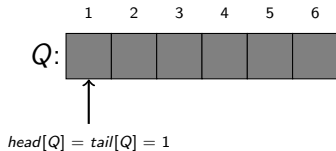
Queue ADT (Cont.)

Auxiliary Operations:

- **FRONT()**: Returns the element at the front without removing it.
- **SIZE()**: Returns the number of elements stored.
- **ISEMPTY()**: Returns a boolean value indicating if the queue is empty or not.

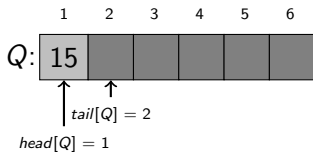


An Integer Array Implementation



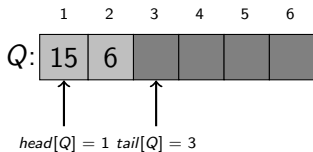
An Integer Array Implementation

ENQUEUE(Q , 15):



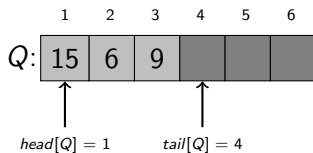
An Integer Array Implementation

ENQUEUE($Q, 6$):



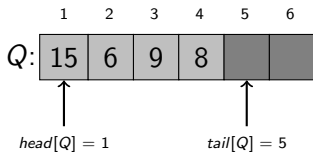
An Integer Array Implementation

ENQUEUE($Q, 9$):



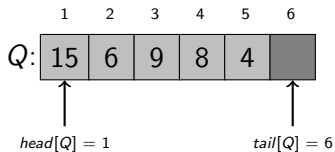
An Integer Array Implementation

ENQUEUE($Q, 8$):



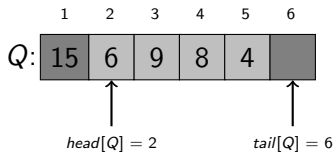
An Integer Array Implementation

ENQUEUE($Q, 4$):



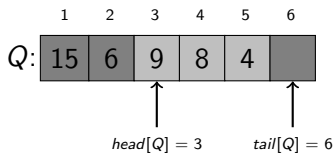
An Integer Array Implementation

DEQUEUE(Q):



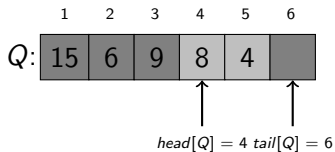
An Integer Array Implementation

DEQUEUE(Q):



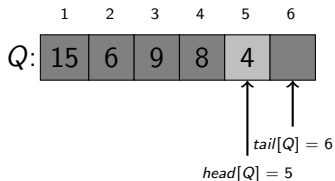
An Integer Array Implementation

DEQUEUE(Q):



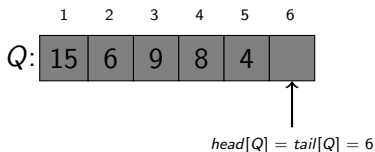
An Integer Array Implementation

DEQUEUE(Q):



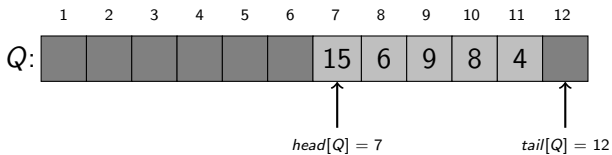
An Integer Array Implementation

DEQUEUE(Q):



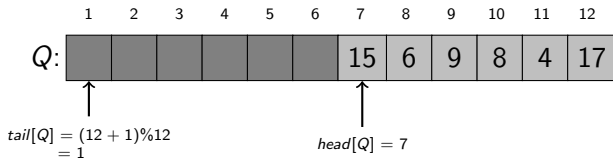
Queue Empty!!

An Improvement



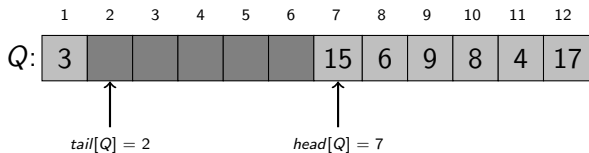
An Improvement

ENQUEUE($Q, 17$):



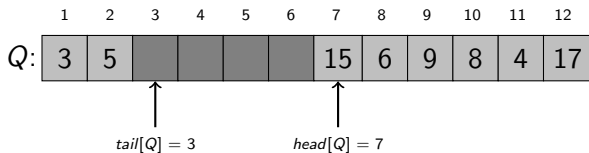
An Improvement

ENQUEUE($Q, 3$):



An Improvement

ENQUEUE($Q, 5$):



An Improvement

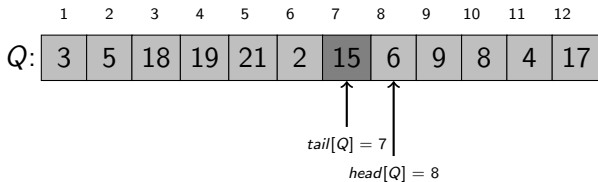
DEQUEUE(Q):



An Improvement

Queue Full: $head[Q] = (tail[Q] + 1) \bmod n$.

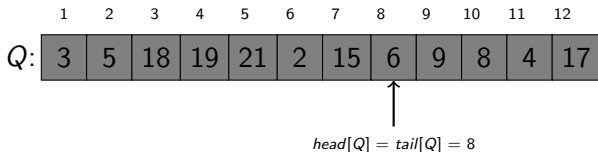
Error: Overflow.



An Improvement

Queue Empty: $head[Q] = tail[Q]$.

Error: Underflow.



ENQUEUE(Q, x)

Begin

If ($head[Q] = (tail[Q] + 1) \bmod n$)
 return “**overflow error**”

$Q[tail[Q]] \leftarrow x$;

If ($tail[Q] = length[Q]$)
 $tail[Q] \leftarrow 1$;

Else

$tail[Q] \leftarrow tail[Q] + 1$;

End

DEQUEUE(Q)

Begin

If ($head[Q] = tail[Q]$)
 return “**underflow error**”

$x \leftarrow Q[head[Q]];$

If ($tail[Q] = length[Q]$)
 $head[Q] \leftarrow 1;$

Else
 $head[Q] \leftarrow head[Q] + 1;$

return $x;$

End

Applications of Queues

- Access to shared resources (e.g., printer).
- Simulations of real world situations of waiting lines (bank teller, flight bookings).
- To efficiently maintain a First-in-first out (FIFO) order on some entities
- In a multitasking operating system, the CPU cannot run all jobs at once, so jobs must be batched up and then scheduled according to order in a queue.
- User input in a game

A C Implementation of a Queue Using An Array

Initialization

```
/* Queue */  
int main() {  
    int head, tail;  
    int Q[len];  
  
    /* Initialisation */  
    head = tail = 0;  
    :  
}
```

ENQUEUE

Insert an element at the tail of the queue Q and redefine tail:

```
/* Enqueue */
int Enqueue(int data, int *Q) {
    /* check if queue is full or not */
    if (head == (tail + 1) % length) {
        printf("\n ERROR: Queue is full\n");
        return FLAG;
    }
    /* insert element at the tail */
    else {
        Q[tail] = data;
        tail = (tail + 1) % length;
    }
    return 0;
}
```

DEQUEUE

Delete and return the element pointed by head of the queue:

```
/* Dequeue */
int Dequeue(int *Q) {
    int x;

    if (head == tail) {    // if queue is empty
        printf("\n ERROR: Queue is empty\n");
        return FLAG;
    }
    /* delete element from the head */
    else {
        x = Q[head];
        head = (head + 1)% length;
    }
    return x;
}
```

FRONT

Return the front element from the queue (if queue is not empty) but do not remove it.

```
/* prints the head of the queue */  
void Front() {  
    if (head == tail) {  
        printf("\n Q is Empty\n");  
        return FLAG;  
    }  
  
    printf("\n Front Element is: %d", Q[head]);  
    return 0;  
}
```

Exercise

Describe the output and final structure of the queue after the following operations:

- ENQUEUE(8)
- ENQUEUE(3)
- DEQUEUE()
- ENQUEUE(2)
- ENQUEUE(5)
- DEQUEUE()
- DEQUEUE()
- ENQUEUE(9)
- ENQUEUE(1)

Linked List

Problems With Arrays

- Many applications require that the number of elements **changes dynamically** as the algorithm progresses.
- **One possibility:** Define all the elements as **arrays** (or **records**) large enough to ensure sufficient storage space.
 - Often a good solution.
 - But requires storage according to the worst case (which may be unknown) which may be *inefficient*.
 - **Reason:** Array size has to be fixed in the beginning.

Problems With Arrays

- May need to perform insertions and deletions in the **middle**.
 - Need to **shift** all other elements.
 - Therefore very costly for large arrays - $\mathcal{O}(n)$.
- This problem is inherent to consecutive representation of arrays (or records).
- For these cases we need **dynamic data structures**.

Handling Lists

- Consider a list of integers

$$\{16, 8, 10, 2, 34, 20, 12, 32, 18, 9, 3\}$$

- It can be thought of as an element 16 followed by another list

$$16 - - \{8, 10, 2, 34, 20, 12, 32, 18, 9, 3\}$$

- Next list can also be thought of as 8 followed by a list

$$16 - - 8 - - \{10, 2, 34, 20, 12, 32, 18, 9, 3\}$$

- \therefore any list can be thought of as an element followed by a list.
- This gives a recursive definition of a [linked list](#).

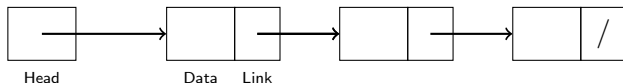
Linked Lists

- Linked lists are the **simplest form** of dynamic data structures.
- The objects are arranged in a **linear order**.
- Provides a simple, flexible representation for dynamic sets, supporting (though not necessarily efficiently) **all** the operations of a dynamic set.

List stored in an array $A[]$

Memory Address	Array Index	List Contents
3200	$A[0]$	36
3204	$A[1]$	42
3208	$A[2]$	20
3212	$A[3]$	16
3216	$A[4]$	38
3220	$A[5]$	40
3224	$A[6]$	12
3228	$A[7]$	54
3232	$A[8]$	82

Linked List



- A **linked list** is a series of connected nodes.
- Each node contains at least
 - A piece of data (any type)
 - Link to the next node in the list
- **Head**: points to the first node
- Links are generated by system.
- The last node points to **nil**.

List Items Stored In A Linked List

Memory Address	Data Contents	Link Contents
2020	36	450
450	42	3600
3600	20	4200
4200	16	4231
4231	38	760
760	40	5555
5555	12	nil

Defining a Node in C

```
typedef struct Node {  
    int nData;  
    struct Node *pNext;  
} Node;  
  
int main() {  
    Node Node1, *pNode2;  
    :
```

- ① Chapter 10.1 & 10.2 of *Introduction to Algorithms* by Thomas H Cormen, Charles E Leiserson, Ronald L Rivest, Clifford Stein.

Thank You for your kind attention!

Questions!!