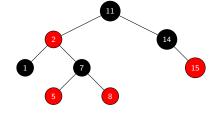
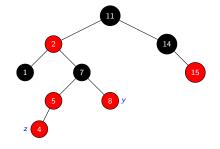
Red-Black Tree: Deletion

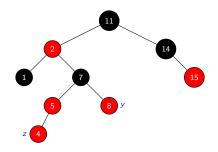
Subhabrata Samajder



IIIT, Delhi Summer Semester, 27th June, 2022 **RBT**: Insertion

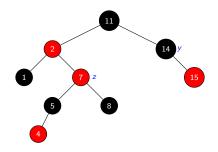






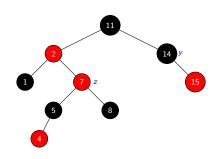
The following cases may arise:

• Case 1: z's uncle is red.



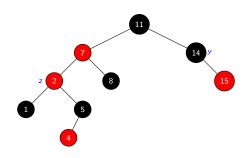
The following cases may arise:

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The following cases may arise:

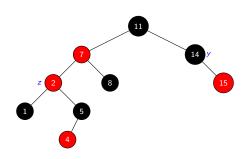
- Case 1: z's uncle is red.
- Case 2: z's uncle is black and z is a right child.



The following cases may arise:

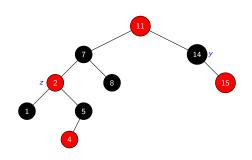
- Case 1: z's uncle is red.
- Case 2: z's uncle is black and z is a right child.

RB-Insert-Fixup(T, z) (Recap)



The following cases may arise:

- Case 1: z's uncle is red.
- Case 2: z's uncle is black and z is a right child.
- Case 3: z's uncle is black and z is a left child.



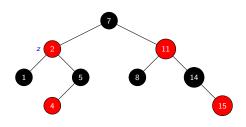
The following cases may arise:

- Case 1: z's uncle is red.
- Case 2: z's uncle is black and z is a right child.
- Case 3: z's uncle is black and z is a left child.

Change Colours:

- Properties 2 is violated.
- Properties 4 is violated for the node 11.

RB-Insert-Fixup(T, z) (Recap)



The following cases may arise:

- Case 1: z's uncle is red.
- Case 2: z's uncle is black and z is a right child.
- Case 3: z's uncle is black and z is a left child.

Change Colours:

- Properties 2 is violated.
- Properties 4 is violated for the node 11.

Follows from the following three part loop-invariant.

Node z is red.

② If p[z] = root[T], then color[p[z]] = BLACK.

- There is at most one violation of either property 2 or 4.
 - Property 2 violation: z is the root and is red.
 - Property 4 violation: colour[z] = colour[p[z]] = RED.

Initialization:

• Recall that z is the red node that was added.

- If p[z] = root[T], then
 - p[z] started out as black and
 - did not change prior to the call of RB-INSERT-FIXUP.

 \bullet Recall that properties 1, 3, and 5 holds when $RB\mbox{-}Insert-Fixup$ is called.

Initialization:

- Property 2 violation:
 - Root must be the newly added node.
 - Since the parent and both the children of z is black ⇒ Property 4 is not simultaneously violated.
 - Implies property 2 is the only violation.

Initialization:

Property 2 violation:

- Root must be the newly added node.
- Since the parent and both the children of z is black ⇒ Property 4 is not simultaneously violated.
- Implies property 2 is the only violation.

Property 4 violation:

- Note that the children of z are black.
- Also, the tree had no prior violations before z was added.
- Implies that colour[z] = colour[p[z]] = RED.
- Also, $root \neq p[z]$ as colour[p[z]] = RED.

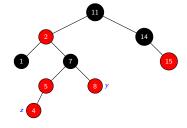
Termination:

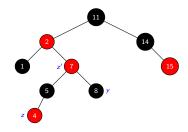
• At termination p[z] = BLACK.

• Thus property 4 is not violated.

• Only violation can be to property 2.

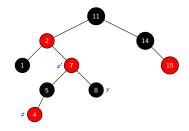
• But Line 16 rectifies it.





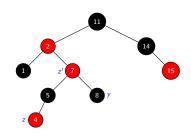
Case 1: z's uncle is red.

- Let z' = p[p[z]] denote the value of z in the next iteration.
- colour[z'] = RED at the start of next iteration.
- Colour of p[z'] = p[p[p[z]]] does not change.
- If p[z'] = root, then the colour of root was black and remains so.



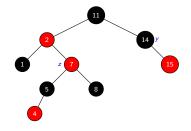
Case 1: z's uncle is red.

- If z' = root, then Case 1 corrected the only violation of property 4.
- colour[z'] = RED and $z' = root \Rightarrow$ property 2 is the only violation.



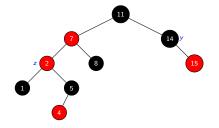
Case 1: z's uncle is red.

- If $z' \neq root$, then Case 1 has not corrected any violation of property 2.
- It corrected the lone violation of property 4.
- It then made colour[z'] = RED and left p[z'] alone.
- If colour[p[z']] = BLACK, there is no violation of property 4.
- If colour[p[z']] = RED, coloring z' red creates one violation of property 4 between z' and p[z'].



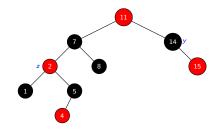
Case 2: z's uncle is black and z is a right child.

- Makes z point to p[z], which is red.
- No further change to z or its color occurs in cases 2 and 3.



Case 2: z's uncle is black and z is a right child.

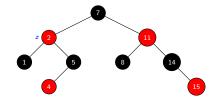
- Makes z point to p[z], which is red.
- No further change to z or its color occurs in cases 2 and 3.



Case 2: z's uncle is black and z is a right child.

Case 3: z's uncle is black and z is a left child.

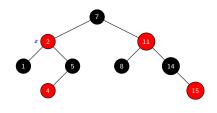
- Makes p[z] black.
- If p[z] = root at the start of the next iteration, then it is black.



Case 2: z's uncle is black and z is a right child.

Case 3: z's uncle is black and z is a left child.

- Makes p[z] black.
- If p[z] = root at the start of the next iteration, then it is black.

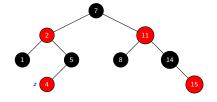


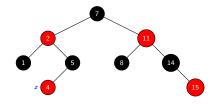
Case 2: z's uncle is black and z is a right child.

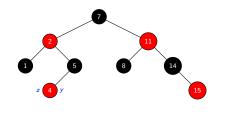
Case 3: z's uncle is black and z is a left child.

- As in Case 1, properties 1, 3, and 5 are maintained in Cases 2 and 3.
- Since node z ≠ root in cases 2 and 3
 ⇒ no violation of property 2.
- Also, Cases 2 and 3 do not introduce a violation of property 2, since the only node that is made red becomes a child of a black node by the rotation in case 3.
- Cases 2 and 3 correct the lone violation of property 4, and they do not introduce another violation.

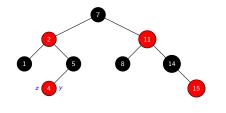
RBT: Deletion



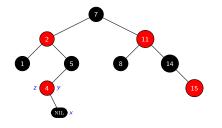




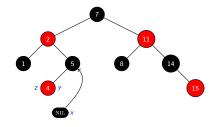
- $2 y \leftarrow z$



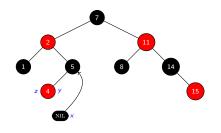
- **1** if left[z] = nil[T] or right[z] = nil[T]
- $y \leftarrow z$



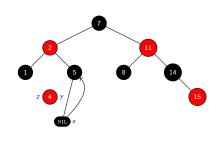
- $2 y \leftarrow z$
- else



- ① if left[z] = nil[T] or right[z] = nil[T]
- $2 y \leftarrow z$
- **3** if $left[y] \neq nil[T]$
- else



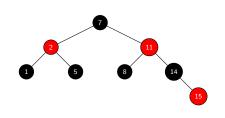
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- ① if left[z] = nil[T] or right[z] = nil[T]
- 2 $y \leftarrow z$
- \bigcirc if $left[y] \neq nil[T]$
- else

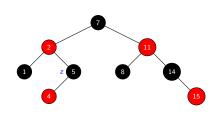
- of p[y] = nil[T]
- 8 else if y = left[p[y]]
- else

3/4



- ① if left[z] = nil[T] or right[z] = nil[T]
- 2 $y \leftarrow z$
- else

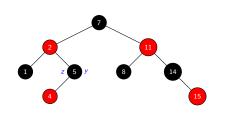
- **8** else if y = left[p[y]]
- else



- ① if left[z] = nil[T] or right[z] = nil[T]

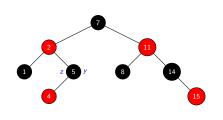
- else

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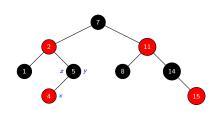
- **1** if left[z] = nil[T] or right[z] = nil[T]
- 2 $y \leftarrow z$
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- else

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- ① if left[z] = nil[T] or right[z] = nil[T]
- 2 $y \leftarrow z$
- **3** if $left[y] \neq nil[T]$
- else

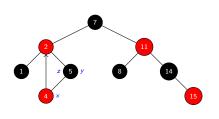
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- **8** else if y = left[p[y]]
- else



- 2 $y \leftarrow z$

- else

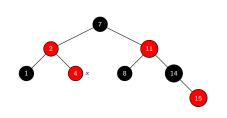
- else



- 2 $y \leftarrow z$

- else

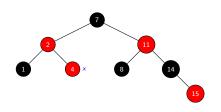
- **9** else if y = left[p[y]]
- else



- $2 y \leftarrow z$

- else

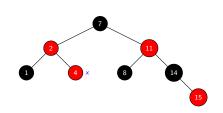
- else



- ① if left[z] = nil[T] or right[z] = nil[T]
- $y \leftarrow z$

- else

- else
- $p[y] \leftarrow x$

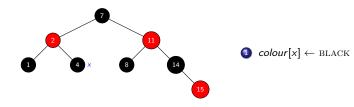


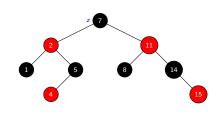
- ① if left[z] = nil[T] or right[z] = nil[T]
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- else

- else
- $p[y] \leftarrow x$
- \square RB-Delete-Fixup(T, x)

RB-DELETE-FIXUP(T, x)

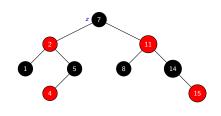




- ① if left[z] = nil[T] or right[z] = nil[T]
- $y \leftarrow z$

- else

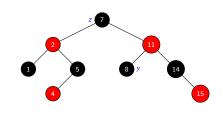
- else
- $p[y] \leftarrow x$
- \mathbb{B} RB-Delete-Fixup(T, x)



- 1 if left[z] = nil[T] or right[z] = nil[T]

- else

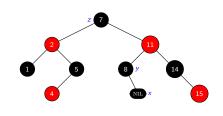
- else
- $p[y] \leftarrow x$
- \mathbb{B} RB-Delete-Fixup(T, x)



- **1** if left[z] = nil[T] or right[z] = nil[T]
- 2 $y \leftarrow z$
- else
- $y \leftarrow \text{Tree-Successor}(z)$

- else

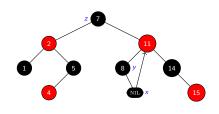
- else
- $p[y]] \leftarrow x$
- if colour[y] = BLACK
- \mathbb{R} B-Delete-Fixup(T, x)



- **1** if left[z] = nil[T] or right[z] = nil[T]
- $2 y \leftarrow z$
- else
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- else

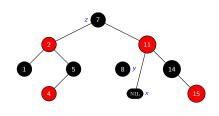
- 📵 else
- **15**if colour[y] = BLACK
- \mathbb{G} RB-Delete-Fixup(T, x)



- **1** if left[z] = nil[T] or right[z] = nil[T]
- 2 $y \leftarrow z$
- else
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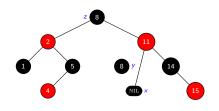
- else

- else
- **15**if colour[y] = BLACK
- \mathbb{G} RB-Delete-Fixup(T, x)



- **1** if left[z] = nil[T] or right[z] = nil[T]
- 2 $y \leftarrow z$
- else
- $y \leftarrow \text{Tree-Successor}(z)$
- of $[y] \neq nil[T]$
- else

- else
- $p[y] \leftarrow x$
- if colour[y] = BLACK
- \mathbb{G} RB-Delete-Fixup(T, x)

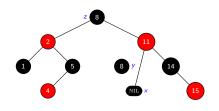


```
① if left[z] = nil[T] or right[z] =
   nil[T]
   y \leftarrow z
else

  y ← Tree-Successor(z)

5 if left[y] \neq nil[T]
0 \quad x \leftarrow left[y]
else
\bigcirc else if y = left[p[y]]
else
  right[p[y]] \leftarrow x
\bigcirc if y \neq z
key[z] \leftarrow key[y]
 \mathbf{0}  if colour[y] = BLACK 
    RB-Delete-Fixup(T, x)
```

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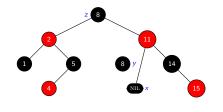


```
① if left[z] = nil[T] or right[z] =
   nil[T]
   y \leftarrow z
else

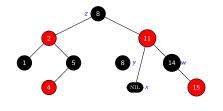
  y ← Tree-Successor(z)

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else
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15 if y \neq z
66 \quad key[z] \leftarrow key[y]
 \mathbf{II}  if colour[y] = BLACK 
    RB-Delete-Fixup(T, x)
```

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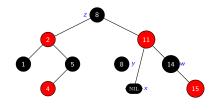
- **1** while $x \neq root[T]$ and colour[x] = BLACK
- \bigcirc colour[x] \leftarrow BLACK



- **1** while $x \neq root[T]$ and colour[x] = BLACK

- \bigcirc colour[x] ← BLACK

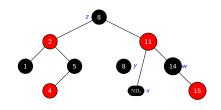
RB-DELETE-FIXUP(T, x)



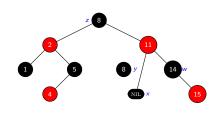
- **1** while $x \neq root[T]$ and colour[x] = BLACK

- if colour[w] = REDCase 1
- **5**colour[x] ← BLACK

RB-DELETE-FIXUP(T, x)



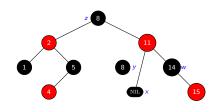
- **1** while $x \neq root[T]$ and colour[x] = BLACK
- - $w \leftarrow right[p[x]]$
- if colour[w] = REDCase 1
- if colour[left[w]] = BLACK and colour[right[w]] = BLACKCase 2
- **6** colour[x] ← BLACK



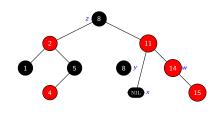
- **1** while $x \neq root[T]$ and colour[x] = BLACK
- if x = left[p[x]]
- - if colour[w] = REDCase 1
- if colour[left[w]] = BLACK and colour[right[w]] = BLACK Case 2
- 6 else if colour[right[w]] =
 BLACK

Case 3

 \bigcirc colour[x] \leftarrow BLACK



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- - $w \leftarrow right[p[x]]$
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- if colour[left[w]] = BLACK and
 colour[right[w]] = BLACK
 Case 2
- 6 else if colour[right[w]] =
 BLACK
 - Case 3
- else if colour[right[w]] = RED
 // Case 4

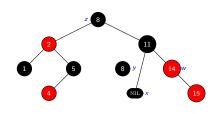


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 colour[right[w]] = BLACK
 Case 2
 - else if colour[right[w]] =
 BLACK

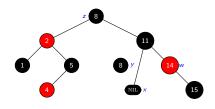
Case 3

- else if colour[right[w]] = RED
 // Case 4
- \bigcirc colour[x] \leftarrow BLACK

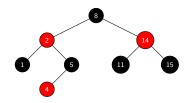


- **1** while $x \neq root[T]$ and colour[x] = BLACK

- if colour[w] = RED
 Case 1
- if colour[left[w]] = BLACK and
 colour[right[w]] = BLACK
 Case 2
- 6 else if colour[right[w]] = BLACK
 - Case 3
- else if colour[right[w]] = RED
 // Case 4
- - $colour[p[x]] \leftarrow BLACK$
- \bigcirc colour[x] \leftarrow BLACK

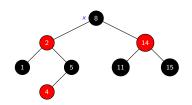


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- if colour[left[w]] = BLACK and
 colour[right[w]] = BLACK
 Case 2
- 6 else if colour[right[w]] =
 BLACK
 - Case 3
- else if colour[right[w]] = RED
 // Case 4
- - $colour[right[w]] \leftarrow BLACK$
- \bigcirc colour[x] \leftarrow BLACK



- **1** while $x \neq root[T]$ and colour[x] = BLACK

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- if colour[left[w]] = BLACK and
 colour[right[w]] = BLACK
 Case 2
 - else if colour[right[w]] = BLACK
 - Case 3
- else if colour[right[w]] = RED
 // Case 4
- $colour[p[x]] \leftarrow BLACK$
- \bigcirc colour[right[w]] \leftarrow BLACK
- \bigcirc colour[x] \leftarrow BLACK



```
1 while x \neq root[T] and colour[x] = BLACK
```

if x = left[p[x]]

 $w \leftarrow right[p[x]]$

if colour[w] = RED
Case 1

if colour[left[w]] = BLACK and
colour[right[w]] = BLACK
Case 2

6 else if colour[right[w]] =
 BLACK

Case 3

else if colour[right[w]] = RED
// Case 4

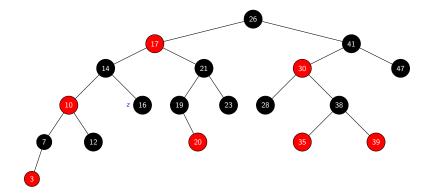
 $oldsymbol{0}$ $colour[p[x]] \leftarrow BLACK$

 \bigcirc colour[right[w]] \leftarrow BLACK

 $x \leftarrow root[T]$

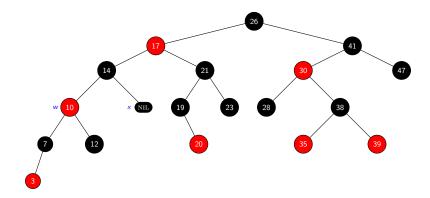
 \bigcirc colour[x] ← BLACK

Delete z = 16.



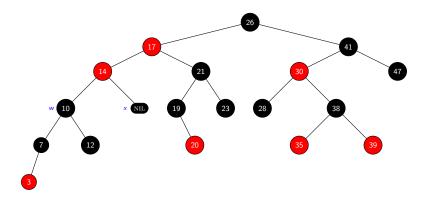
RB-DELETE-FIXUP(T, x)

Case 1: colour[x] = BLACK, x = right[p[x]] and colour[w] = RED (Mirror case!)



Case 1: colour[x] = BLACK, x = right[p[x]] and colour[w] = RED (Mirror case!)

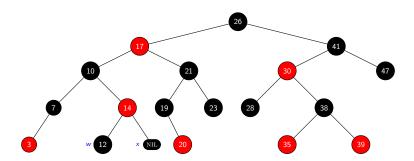
- $colour[w] \leftarrow BLACK$
- $colour[p[x]] \leftarrow RED$



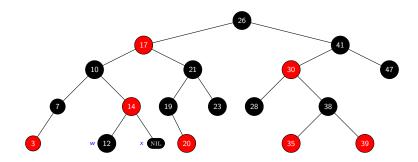
RB-DELETE-FIXUP(T, x)

Case 1: colour[x] = BLACK, x = right[p[x]] and colour[w] = RED (Mirror case!)

- $colour[w] \leftarrow BLACK$
- $colour[p[x]] \leftarrow \text{RED}$
- RIGHT-ROTATE(T, p[x])
- $w \leftarrow left[p[x]]$



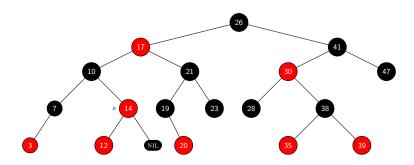
Case 2: colour[left[w]] = BLACK and colour[right[w]] = BLACK



RB-DELETE-FIXUP(T, x)

Case 2: colour[left[w]] = BLACK and colour[right[w]] = BLACK

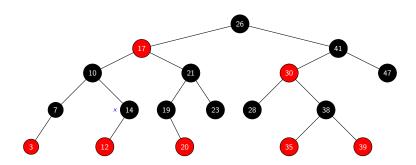
- $colour[w] \leftarrow RED$
- \bullet $x \leftarrow p[x]$



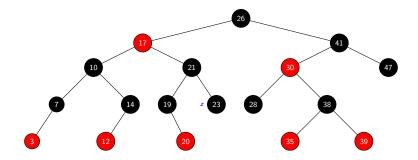
Case 2: colour[left[w]] = BLACK and colour[right[w]] = BLACK

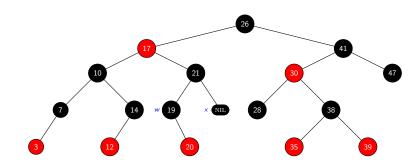
- $colour[w] \leftarrow \text{RED}$
- \bullet $x \leftarrow p[x]$

 $colour[x] \leftarrow BLACK$ (outside the while loop)



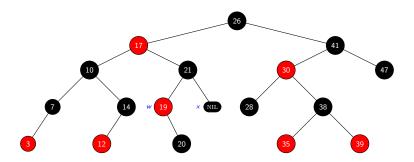
Delete z = 23.





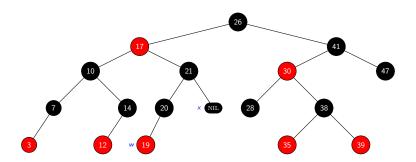
RB-DELETE-FIXUP(T, x)

- $colour[right[w]] \leftarrow BLACK$
- $colour[w] \leftarrow RED$

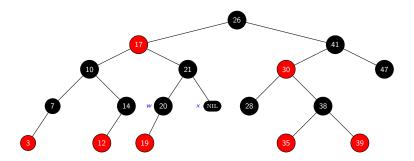


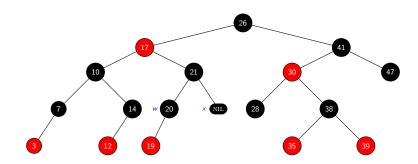
RB-DELETE-FIXUP(T, x)

- $colour[right[w]] \leftarrow BLACK$
- $colour[w] \leftarrow RED$
- Left-Rotate(T, w)

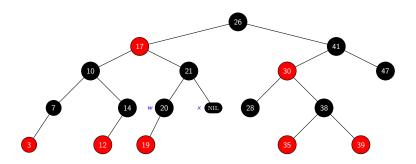


- $colour[right[w]] \leftarrow BLACK$
- $colour[w] \leftarrow RED$
- Left-Rotate(T, w)
- $w \leftarrow left[p[x]]$

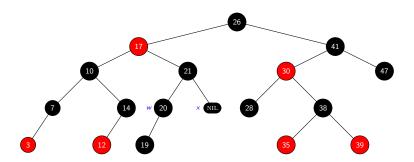




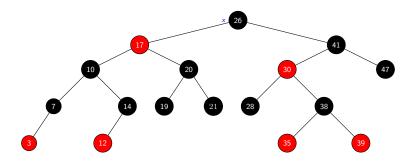
- $colour[w] \leftarrow colour[p[x]]$
- $colour[p[x]] \leftarrow BLACK$



- $colour[w] \leftarrow colour[p[x]]$
- $colour[p[x]] \leftarrow BLACK$
- $colour[left[w]] \leftarrow BLACK$



- $colour[w] \leftarrow colour[p[x]]$
- $colour[p[x]] \leftarrow BLACK$
- $colour[left[w]] \leftarrow BLACK$
- RIGHT-ROTATE(T, p[x])
- $x \leftarrow root[T]$



Thank You for your kind attention!

Books and Other Materials Consulted

• Red-Black Tree part taken from Chapter 13 of the Introduction to Algorithms book by Thomas H Cormen, Charles E Leiserson, Ronald L Rivest, Clifford Stein.

Questions!!