

Introduction to Data Structures

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Definition

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More precisely, a **data structure** is a collection of data values, the relationships among them, and the functions or operations that can be applied to the data.

- They are the building blocks of computer algorithms.
- Design of an algorithm must be based on a thorough understanding of data structure techniques and costs.

Abstract Data Type (ADT)

- It is a useful notion in the study of data structures.
- Normally, when we write a program, we have to specify the data type (e.g., **integers**, **reals**, **characters**).
 - Also called **primitive data structures**.
 - Directly operated upon by the machine-level instructions.

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- **Condition:** First-in first-out (FIFO).

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Example:

- **Frequent Operations:** Insertions and deletions.
 - **Condition:** First-in first-out (FIFO).
 - **Data Structure:** Queue.
-
- It is more convenient and more general to design algorithms for these operations without specifying the data type of the items.

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- If our needs matches the definition of an ADT, we use it.
- Make the algorithm-design process **more modular**.

Abstract Data Type: Definition

Definition

An **abstract data type (ADT)** is a **mathematical model** for data types, where a data type is defined by its **behavior (semantics)** from the point of view of a **user** of the data, specifically in terms of possible **values**, possible **operations** on data of this type, and the **behavior** of these operations.

In contrast a **Data structures** requires

- concrete representations of data, and
- are of the point of view of an **implementer** and **not** a **user**.

ADT Example: Integers

Integers are an ADT, defined by the

- **Values:** $\dots, -2, -1, 0, 1, 2, \dots$, and by the
- **Operations:** '+', '-', '.', '/', '<', etc.
- **Behavior:**
 - Obeying various axioms (associativity and commutativity of addition etc.).
 - Preconditions on operations (cannot divide by zero).
 - Must be independent of how the integers are represented.
- **Representation:**
 - Typically represented in 2's complement.
 - Can also be binary-coded decimal or in 1' complement.
- **User:**
 - Abstracted from the concrete choice of representation.
 - Simply use it as data types.

Elementary Data Structures

- **Sets** are also fundamental to computer science.
- **Dynamic Sets:** Sets manipulated by algorithms can *grow*, *shrink*, or *change over time*.
- Algorithms may require several different types of operations to be performed on sets.
- **Dictionary:** A dynamic set that supports *insertion*, *deletion*, and *membership testing*.
- The best way to implement a dynamic set depends upon the operations that must be supported.

- **Elements:** Generic name for an unspecified data type.

Example: Integer, a set of integers, a text file, etc.

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- **Assumptions:**

- ① Elements can be compared for equality.
- ② Elements are taken from a **totally ordered set**.
- ③ Elements can be copied.

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- **Assumptions:**

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- ③ Elements can be copied.

- **Assumption:** All these operations take **unit amount of time**.

Keys:

- An identifying field of the objects in a dynamic set.
- If all the keys are **different**, then we think of the dynamic set as being a **set of key values**.

Satellite Data:

- Data carried around in other object fields but are otherwise not part of the implementation.

Operations on a Dynamic Set

- Can be grouped into two categories:
 - **Queries:** Which simply return information about the set.
 - **Modifying Operations:** Which change the set.

Modifying Operations

INSERT(S, x): Augments the set S with the element pointed to by x .

DELETE(S, x): Given a pointer to an element x in the set S , removes x from S .

Note: These operations uses a pointer to x and not a key value.

Queries

SEARCH(S, k): Given a set S and a key value k , returns

- a pointer x to an element in S such that $\text{key}[x] = k$, or
- **nil** if no such element belongs to S .

MINIMUM(S): Returns a pointer to the element with smallest key.

MAXIMUM(S): Returns a pointer to the element with largest key.

SUCCESSOR(S, x): Given an element x , returns

- a pointer to the next larger element in S , or
- **nil** if x is the maximum element.

PREDECESSOR(S, x): Given an element x , returns

- a pointer to the next smaller element in S , or
- **nil** if x is the minimum element.

ADT: Arrays

Arrays

- Row of elements of the same type.
- The **size** of an array is the number of elements in that array.
- The **size** must be **fixed**.
- \therefore all the elements are of the same type.
- Thus amount of memory is known **a priori**.
- Every element of an array can be accessed in **constant time**.

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- Thus amount of memory is known **a priori**.
- Every element of an array can be accessed in **constant time**.
- **Drawbacks:**
 - Cannot be used to store elements of different types (or sizes).
 - The size of an array cannot be changed dynamically.

ADT: Records

Records

- Similar to arrays, except elements can be of different types.
- A **record** is thus a list of elements of different types.
- The exact combination of types is **fixed**.
- \therefore the **storage size** of a record is known in advance.
- Each element in a record can be accessed in **constant time**.

Records

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- A **record** is thus a list of elements of different types.
- The exact combination of types is **fixed**.
- \therefore the **storage size** of a record is known in advance.
- Each element in a record can be accessed in **constant time**.
 - This is accomplished by keeping an **array**.
 - Elements are then accessed by consulting the array.
 - The exact program that maintains the array is created automatically by the **compiler**.

Records: An Example

record example1

Begin

Int1 : integer;

Int2 : integer;

Ar1 : array[1 ... 20] of integer;

Ar2 : array[1 ... 20] of integer;

Ar3 : array[1 ... 20] of integer;

Int3 : integer;

Int4 : integer;

Int5 : integer;

Int6 : integer;

Name1 : array[1 ... 11] of character;

Name2 : array[1 ... 12] of character;

End

- The array contains starting relative locations of all the elements.
- Thus 'Int6' starts at byte number 261 ($= 2 \cdot 4 + 3 \cdot 20 \cdot 4 + 3 \cdot 4 + 1$)
- Like arrays, the storage for a record is always **consecutive**.

Records: An Example

```
record example1
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```
Begin
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  Int1 : integer;
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  Ar1 : array[1 ... 20] of integer;
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  Ar2 : array[1 ... 20] of integer;
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  Ar3 : array[1 ... 20] of integer;
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  Int3 : integer;
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  Int4 : integer;
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```
  Int5 : integer;
```

```
  Int6 : integer;
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```
  Name1 : array[1 ... 11] of character;
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  Name2 : array[1 ... 12] of character;
```

```
End
```

- The array contains starting relative locations of all the elements.
- Thus 'Int6' starts at byte number 261 ($= 2 \cdot 4 + 3 \cdot 20 \cdot 4 + 3 \cdot 4 + 1$)
- Like arrays, the storage for a record is always **consecutive**.
- **Drawback:** It is **not** possible to add elements dynamically.

Records in C

```
struct Record {  
    int Int1;  
    int Int2;  
    int Ar1[20];  
    int Ar2[20];  
    int Ar3[20];  
    int Int3;  
    int Int4;  
    int Int5;  
    int Int6;  
    char Name1[11];  
    char Name2[12];  
} Example1;
```

Records in C

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struct Record {  
    int Int1;  
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    int Ar1[20];  
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    int Int3;  
    int Int4;  
    int Int5;  
    int Int6;  
    char Name1[11];  
    char Name2[12];  
} Example1;
```

sizeof(Example1)?

Books Consulted

- ① Chapter 10 of *Introduction to Algorithms* by Thomas H Cormen, Charles E Leiserson, Ronald L Rivest, Clifford Stein.

Thank You for your kind attention!

Questions!!