CSE 358 - ASSIGNMENT 8

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Question 1

q1.l:

```
≣ q1.l
     You, 1 second ago | 1 author (You)
     %{
     #include "q1.tab.h"
     void yyerror(char *);
     %}
     %%
     [0-9]+
                    { yylval.num = atoi(yytext); return NUMBER; }
     [+\-*/()]
                    { return yytext[0]; }
                    ; // Ignore only spaces/tabs
     [\t]
                     { yyerror("Invalid character"); }
11
     %%
13
     int yywrap() { return 1; }
```

q1.y:

```
≡ q1.y
    %{
    #include <stdio.h>
    #include <stdlib.h>
     int yylex(void);
    void yyerror(char *);
    %}
    %union { int num; }
    %token <num> NUMBER
    %type <num> expr
    %left '+' '-'
    %left '*' '/'
    %right UMINUS
    %%
    input:
             /* empty */
            | input expr '\n' { printf("\n"); }
                           { printf("%d ", $1); }
     expr:
            NUMBER
21
             '(' expr ')' { $$ = $2; }
            | expr '+' expr { printf("+ "); }
            | expr '/' expr { printf("/ "); }
              '-' expr %prec UMINUS { printf("- "); $$ = $2; }
    %%
     void yyerror(char *s) {
        fprintf(stderr, "Error: %s\n", s);
     int main() {
        printf("Enter expression: ");
        fflush(stdout); // Force prompt display
        yyparse();
        return 0;
```

Output:

```
arnav@arnav-IdeaPad-Gaming-3-15ACH6:~/Desktop/Compiller-Techniques/LAB 8$ ./q1
Enter expression: 2+3/5
2 3 5 / +
```

Question 2

q2.1

```
≣ q2.l
    %{
     #include "q2.tab.h"
    void yyerror(char *); // Explicit error declaration
     extern int count;  // Access YACC's counter
     %}
     %%
                    { yylval.digit = 0; return BINARY_DIGIT; }
                    { yylval.digit = 1; return BINARY DIGIT; }
                    { return DOT; }
                    ; // Ignore whitespace
11
     [ \t\n]
                    { yyerror("Invalid binary character"); }
12
13
     %%
    int yywrap() { return 1; }
```

Q2.y

```
≡ q2.y
     #include <stdio.h>
     #include <math.h>
     void yyerror(char *);
     int yylex(void);
     int count = 0; // Track fractional digit position
     %}
     %union {
11
        int digit;
         double val;
     %token <digit> BINARY DIGIT
     %token DOT
     %type <val> number integer part fractional part
     %%
     number:
                integer part DOT fractional part {
                     $$ = $1 + $3;
                     printf("Decimal value: %.4f\n", $$);
     integer part: /* empty */ { $$ = 0.0; }
                 | integer part BINARY DIGIT { $$ = $1 * 2 + $2; }
     fractional part: /* empty */ { $$ = 0.0; count = 0; }
                 | fractional part BINARY DIGIT {
                     count++;
                     $$ = $1 + $2 * pow(2, -count);
     %%
     void yyerror(char *s) {
         fprintf(stderr, "Error: %s\n", s);
     int main() {
         yyparse();
         return 0;
```

Output

```
■ arnav@arnav-IdeaPad-Gaming-3-15ACH6:~/Desktop/Compiller-Techniques/LAB 8$ ./q2
111.10
Decimal value: 7.5000

$\frac{1}{2}\arnav@arnav-IdeaPad-Gaming-3-15ACH6:~/Desktop/Compiller-Techniques/LAB 8$
```

Question 3

Question 3: Error Handling and Recovery in Compiler Design

Error Handling

Error Handling in compiler design identifies and reports errors in source code while maintaining the compilation process. The errors include:

- i) lexical Errors: incorrect grammar (eg missing semicolon like int x = 10)
- ii) Senantic Errors: Meaning errors (eg. undéclared variable y used)
- iii) Runtime Errors: Errors during program execution (eg. division by zero)
- iv) Linking Errors: Issues during code linking (eg missing library)

Error Recovery

methods for recovery after detecting errors:

- i) Panic Mode: Discards input until a synchronizing token is found (ie skipping to the next semicolon)
- ii) Phrase level: Tries to replace or insert tokens to continue passing (eg changing { to)).
- iii) Error Productions: Allows predefined error handling within grammar rules.
- iv) Contextual Recovery: Analyse surrounding code to decide how to recover
- v) Back tracking: Tries alternate parsing strategies when one path fails.

Question 4

Question 4: Detailed Under Standing of Compiler Design and Phase A compiler translates high level source code into machine code or intermediate code to execute the program Compiler Phases:

i) Lexical Analysis: input > source code output >> Tokens (eq. keywords, operators)

ii) Syntax Analysis: input > Tokens
output > Parse Tree (structure showing program's
Syntax)

(ii) Semantic Analysis: input -> Parse Tree

Output -> Intermediate Code representation

(eg: Three-order Code)

iv) Optimization: Input > Intermediate Code
Output > Optimized Intermediate Code

V) Code Generation: Input > Opinised Intermediate Code
Output > Machine Assembly code

vi) Code Optimizer: Input -> Machine Code (optional) Output -> Machine Optimized Code

vii) Code Emission Input -> Final Code and Linking : Output -> Executable file.

GitHub

https://github.com/arnavjain2710/Compiller-Techniques/tree/main/LAB%208