

1. Let F be a field, $A \in M_{n \times n}(F)$ and $\lambda \in F$. Prove that for all $\lambda \in F$, the sets $\{v \in F^n | Av = \lambda v\}$ and $\{v \in F^n | A^t v = \lambda v\}$ are subspaces of the same dimension. Conclude that λ is an eigenvalue of A if and only if λ is an eigenvalue of A^t .

Definition 1. A matrix $A \in M_{n \times n}(\mathbb{R})$ is *column-stochastic* if all of its entries are non-negative and the sum of the entries in each column is equal to 1.

2. Prove that a column-stochastic matrix has $\lambda = 1$ as an eigenvalue.
3. (a) Let $a_1, \dots, a_n \in \mathbb{R}$. Prove that $|\sum_{i=1}^n a_i| < \sum_{i=1}^n |a_i|$ if and only if there are i, j such that $a_i > 0$ and $a_j < 0$.
 (b) Prove that if $v, w \in \mathbb{R}^n$ are linearly independent, then there is some $u \in \text{span}\{v, w\}$ such with at least one positive entry and one negative entry.
 (c) Let $A \in M_{n \times n}(\mathbb{R})$ be a column-stochastic matrix with all entries positive. Prove that $\{v \in \mathbb{R}^n | Av = v\}$ is a 1-dimensional vector space.
 (d) With A as in (c), conclude there is a unique vector v such that $Av = v$, all of its entries are positive and the sum of its entries is 1.
4. Let $n \in \mathbb{N}$. Let $\mathcal{P} := \{1, \dots, n\}$ (elements of \mathcal{P} are referred to as pages) and let $\mathcal{L} \subseteq \mathcal{P}^2$ be some set of ordered pairs of \mathcal{L} (elements of \mathcal{L} are called links). If $(a, b) \in \mathcal{L}$, we say there is a link from a to b . Fix some $0 < p < 1$. Consider the following random process :

- In stage 0: we pick a page with uniform distribution (i.e., the probability of every element of \mathcal{P} to be picked is $1/n$).
- In stage $k + 1$: Assuming in stage k we picked the page a , in stage $k + 1$, with probability p we pick a new page out of \mathcal{P} with uniform distribution, and with probability $1 - p$ we pick, with uniform distribution, some page out of the pages with a link from a .

Namely:

- If $(a, b) \notin \mathcal{L}$, then the probability of picking b is $p \cdot 1/n$.
- If $(a, b) \in \mathcal{L}$ and there are l many links from a , then the probability of picking b is $p \cdot 1/n + (1 - p) \cdot 1/l$.

Find a matrix $A \in M_{n \times n}(\mathbb{R})$ and a vector $v \in \mathbb{R}^n$ such that the probability of picking page i in stage k is the i -th coordinate of $A^k v$. Prove A is column-stochastic.

In this sheet, you almost developed the basic PageRank, Google's ranking of web pages. For details, see the back of this page.

The process described in Question 4 describes a random surfer on the web. The surfer starts by surfing to a random web page picked at random, then with probability p (where p is quite small, e.g., 0.15) the surfer jumps to a completely random new page and with probability $(1 - p)$ the surfer follows one of the links on the page currently at.

It can be shown that the process Av, A^2v, A^3v, \dots converges on each coordinate. Moreover, the limit vector, call it v^* , is precisely the unique vector found in Question 3d. However, this requires some more analysis. The i -th coordinate of v^* is the PageRank of page i .

Of course, there is more to Google's algorithm, e.g., what's the magic p , giving some links more weight than others etc. More importantly, how to compute v^* ? Having said that, still, these are the basic foundations of Google's PageRank.

As for the question of computing (or at least, approximating) v^* , this can be solved by taking powers of matrices efficiently, a matter that will be addressed later in this module.

For further information, see:

<http://pi.math.cornell.edu/~mec/Winter2009/RalucaRemus/index.html>