CS621/CSL611 Quantum Computing For Computer Scientists

Quantum Architecture

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Classical Gates From Quantum Perspective

How should we represent our logical gates?

- We represent n input bits as a $(2^n \times 1)$ matrix and m output bits as a $(2^m \times 1)$ matrix
- When one multiplies a $(2^m \times 2^n)$ matrix with a $(2^n \times 1)$ matrix, the result is a $(2^m \times 1)$ matrix.

$$(2^m \times 2^n) \star (2^n \times 1) = (2^m \times 1)$$

- Bits → column vectors
- Logic gates → matrices

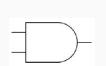
• To study classical gates from the point of view of matrices.

$$\mathsf{NOT} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$\mathsf{NOT}\ket{0} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \ket{1}$$

$$\mathsf{NOT}\ket{1} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \ket{0}$$

AND Classical Gates



$$\mathsf{AND} = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

• AND
$$|11\rangle = |1\rangle$$

$$\begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

• AND $|01\rangle = |0\rangle$

$$\begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

ullet Similarly, AND |10
angle = |0
angle and AND $|00
angle = |0
angle^1$

 $^{^{1}}$ We are **allowed only** to multiply these classical gates with vectors that represent classical states, i.e., column matrices with a single 1 entry and all other entries 0.

• Find the matrix for the OR Gate



• The matrix:

NAND Classical Gates

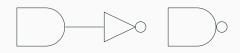


$$NAND = \begin{bmatrix} \mathbf{00} & \mathbf{01} & \mathbf{10} & \mathbf{11} \\ \mathbf{0} & 0 & 0 & 0 & 1 \\ \mathbf{1} & 1 & 1 & 0 \end{bmatrix}$$

Note

Column names correspond to the inputs and the row names correspond to the outputs. 1 in the jth column and ith row means that on entry j the matrix/gate will output i.

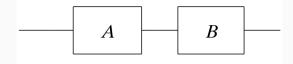
• Another way to determine the NAND gate



$$NOT \star AND = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \star \begin{bmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix} = NAND$$

Multiple Gates

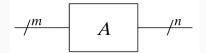
• Time goes from left to right.



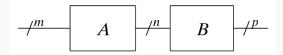
- A computation of A followed by B
- What does this mean in terms of the corresponding matrix multiplications?

Sequential Operations

• What is size of the matrix for A?

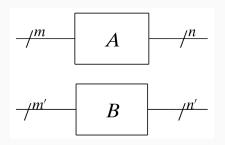


• What are the sizes for A and B?



- Above sequential operation corresponds to $B \star A$
- What is the size of the matrix corresponding to $B \star A$?

Parallel Operations



- $A \rightarrow 2^n \times 2^m$
- $B \rightarrow 2^{n'} \times 2^{m'}$

- Here we have A acting on some bits and B on others
- This is captured by tensor product: A ⊗ B
- $A \otimes B \rightarrow 2^n 2^{n'} \times 2^m 2^{m'}$
- So $A \otimes B$ is of size:

$$2^{n+n'} \times 2^{m+m'}$$

- AKA Kronecker product
- In general, for any two matrices

$$A = \begin{pmatrix} a_{1,1} & a_{1,2} & \cdots & a_{1,m} \\ a_{2,1} & a_{2,2} & \cdots & a_{2,m} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n,1} & a_{n,2} & \cdots & a_{n,m} \end{pmatrix}_{n \times m} B = \begin{pmatrix} b_{1,1} & b_{1,2} & \cdots & b_{1,l} \\ b_{2,1} & b_{2,2} & \cdots & b_{2,l} \\ \vdots & \vdots & \ddots & \vdots \\ b_{k,1} & b_{k,2} & \cdots & b_{k,l} \end{pmatrix}_{k \times l}$$

• $A \otimes B$ is defined by the $nk \times ml$ matrix²

$$A \otimes B = \begin{pmatrix} a_{1,1}B & a_{1,2}B & \cdots & a_{1,m}B \\ a_{2,1}B & a_{2,2}B & \cdots & a_{2,m}B \\ \vdots & \vdots & \ddots & \vdots \\ a_{n,1}B & a_{n,2}B & \cdots & a_{n,m}B \end{pmatrix}_{nk \times ml}$$

²The definition works for vectors by thinking of them as one column matrices.

Tensor Product Properties

• Associative: For any choice of matrices A, B and C:

$$(A \otimes B) \otimes C = A \otimes (B \otimes C)$$

• For any choice of matrices A, B, C and D^3

$$(A \otimes B)(C \otimes D) = (AC) \otimes (BD) \tag{1}$$

- Distributive: $A \otimes (B + C) = A \otimes B + A \otimes C$
- Scalar multiplication⁴: $(\alpha A) \otimes B = A \otimes (\alpha B) = \alpha (A \otimes B)$
- Not-commutative: $A \otimes B \neq B \otimes A$

Eq (1) is of special significance while dealing with quantum circuits

³Juxtaposition implies multiplication

⁴Scalars "float freely" through the tensor product

• Let X and Y be qubits having associated superpositions

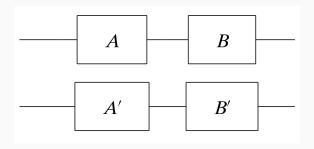
$$v = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} \qquad w = \begin{pmatrix} \gamma \\ \delta \end{pmatrix}$$

- The superposition of the pair (X,Y) is $v\otimes w=\begin{pmatrix} \alpha\gamma\\ \alpha\delta\\ \beta\gamma\\ \beta\delta \end{pmatrix}$
- Not every superposition can be written as a tensor product

$$\begin{pmatrix} \frac{1}{\sqrt{2}} \\ 0 \\ 0 \\ \frac{1}{\sqrt{2}} \end{pmatrix} \leftarrow$$
 Recall: For quantum case, this type of correlation is called an entanglement

Analyzing Eq (1)

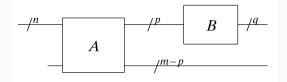
• $(B \otimes B')(A \otimes A') = (BA) \otimes (B'A')$



 To what does this correspond in terms of performing different operations on different (qu)bits?

Operating On Parts

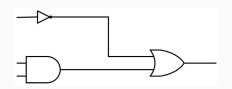
- How to treat the m-p bits that are not operated on by B?
- Here B is a $2^q \times 2^p$ matrix



- Use the $(2^{m-p} \times 2^{m-p})$ Identity matrix I_{m-p}
- The entire circuit can be represented by the following matrix:

$$(B\otimes I_{m-p})A$$

Circuit Example



$$NOT \otimes AND = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$OR \star (NOT \otimes AND) = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$