Assignment 5 — Traps and Interrupts Due: Wednesday, May 11th

In a real computer, there are typically far more running processes than there are physical CPU cores. In such situations, rather than having a process waste resources by polling while it waits for user input, another process can be allowed to run during that time¹.

Deliverables:

GitHub Classroom: https://classroom.github.com/a/prRDtFJH

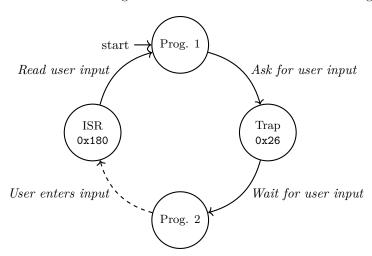
Required Files: trap26.asm

Optional Files: none

Part 1: Time-Sharing

Implementing interrupt-driven I/O requires a few extra steps. Notably, recall that in order to enable keyboard interrupts, bit 14 must be set to a '1' in the KBSR. However, the KBSR is memory-mapped to 0xFE00, which is only accessible in privileged mode.

In order to properly use interrupts, you will need to implement *both* a trap service routine and an interrupt service routine, which will need to work together to transition to and from the calling program:



- 1. Execution begins with Program 1, which reads a single character. Rather than using GETC, Program 1 instead calls Trap 0x26.
- 2. The trap service routine sets up keyboard interrupts, but it does *not* return to Program 1 (which was its caller). Instead, it "returns" to Program 2.
- 3. Program 2 is allowed to execute while Program 1 is waiting for input.

Once the user types a character, an interrupt is triggered; ISR 0x180 responds.

4. The interrupt service routine reads the typed character, but it does *not* return to Program 2 (which was interrupted). Instead, it "returns" to Program 1.

From the perspective of Program 1, Trap 0x26 and ISR 0x180 behave as two halves of a single service routine.

¹Certainly, the mechanism implemented here is a gross oversimplification of how a real operating system implements *time-sharing*. However, within the confines of the LC-3, it will illustrate the fundamentals.

Part 2: Trap 0x26

Complete the service routine labeled TRAP26 in trap26.asm. This service routine is responsible for setting up keyboard interrupts and switching to execution of Program 2:

- Your implementation must place the address of ISR180 (which you will implement shortly) into the keyboard's entry in the Interrupt Vector Table.
- It must enable keyboard interrupts by setting bit 14 of the KBSR to '1'.
- It must "return" to Program 2. When TRAP26 is called, the supervisor stack will contain (left):



...the service routine must swap Program 1's PC and PSR with those of Program 2 (right), which will cause an RTI to return to the latter. Moreover, TRAP26 must store all of Program 1's state:

...the state of Program 2 must then be loaded²:

$$\cdot$$
 PC \cdot R0 \cdot R2 \cdot R4 \cdot R7 \cdot PSR \cdot R1 \cdot R3 \cdot R5

Note that Program 1's R0 need not be saved; it will eventually be used to return a typed character. You may also assume that R6 is used exclusively for the supervisor stack, and should not be touched.

Part 3: ISR 0x180

Complete the service routine labeled ISR180 in trap26.asm. This service routine is responsible for responding to keyboard interrupts and switching back to execution of Program 1:

- · Your implementation must read the typed character into RO. It may not call GETC or use polling.
- It must disable keyboard interrupts by setting bit 14 of the KBSR to '0'.
- It must "return" to Program 1. When ISR180 is called, the supervisor stack will contain (left):



...the service routine must swap Program 2's PC and PSR with those of Program 1 (right), which will cause an RTI to return to the latter. Moreover, ISR180 must store all of Program 2's state:

...the state of Program 1 must then be loaded:



That is to say, ISR180 must restore all of Program 1's state that was saved by TRAP26, and it must save all of Program 2's state that will be restored if TRAP26 is ever called again.

²The initial state of Program 2 has been defined in trap26.asm, so that the first call to TRAP26 will begin Program 2.

Part 4: Testing

Consider the code in program1a.asm and program2.asm. Here, Program 1 uses TRAP x26 to read a single character using interrupt-driven I/O. Program 2 simply infinitely, periodically increments the contents of R1.

Not only must program1a, program2, and trap26 be assembled and loaded into the LC-3, but memory location 0x0026 must be manually set to 0x1000 to enable TRAP x26.

- Throughout this assignment, you may not alter any of the code for Programs 1 or 2.
- Depending on when characters are typed, the state of Program 2 as saved by ISR180 and restored by TRAP26 will vary. It need *not* be possible to rerun Program 1 by manually resetting the PC to 0x3000.

You may assume that Program 2 will neither alter the keyboard's entry in the Interrupt Vector Table nor itself use TRAP x26. You may further assume that the user will wait long enough to type a character that a keyboard interrupt will not be triggered while TRAP26 is running.

For example:

```
[Program 1] Reading...
[Program 2] R1: 0...
[Program 2] R1: 1...
[Program 2] R1: 2...
[Program 1] Read 'a'.
```

...where the number of lines printed will naturally depend on how long the user waits before typing a character.

Your implementation will be tested with other variations of Program 1 and Program 2. For example, consider program1b.asm, which invokes TRAP x26 within an infinite loop:

```
[Program 1] Reading...
[Program 2] R1: 0...
[Program 2] R1: 1...
[Program 1] Read 'f'.
[Program 1] Reading...
[Program 2] R1: 2...
[Program 1] Read 'o'.
[Program 1] Reading...
[Program 2] R1: 3...
[Program 2] R1: 4...
[Program 2] R1: 5...
[Program 1] Reading...
[Program 2] R1: 6...
[Program 2] R1: 6...
```

Part 5: Submission

The following files are required and must be pushed to your GitHub Classroom repository by the deadline:

• trap26.asm — A working assembly trap and ISR for interrupt-driven keyboard input, as specified.

The following files are optional:

· none

Any files other than these will be ignored.