Homework 5

Exercise 1

$$\begin{pmatrix} -a_r & 0 \\ -l_{\Pi+r} & 1 \end{pmatrix} \begin{pmatrix} dr \\ dm \end{pmatrix} = \begin{pmatrix} a_f df - dy \\ l_{\Pi+r} d\Pi + l_y dy \end{pmatrix}$$

$$\begin{array}{ll} b^{-1} = \left(\begin{array}{cc} -\frac{1}{a_r} & 0 \\ -\frac{1_{r+\Pi}}{a_r} & 1 \end{array} \right) \\ \left(\begin{array}{c} dr \\ dm \end{array} \right) = \left(\begin{array}{cc} -\frac{1}{a_r} & 0 \\ -\frac{1_{r+\Pi}}{a_r} & 1 \end{array} \right) \left(\begin{array}{c} a_f \, df - dy \\ 1_{\Pi+r} \, d\Pi + 1_y \, dy \end{array} \right) \\ \left(\begin{array}{c} dr \\ dm \end{array} \right) = \left(\begin{array}{c} \frac{dy - df \, a_f}{a_r} \\ dy \, 1_y + \frac{(dy - df \, a_f + d\Pi \, a_r) \, 1_{r+\Pi}}{a_r} \end{array} \right) \end{array}$$

We can find $\frac{dr}{df}$ by dividing the first row of the matrix by df.

Exercise 2

```
Maximize x y s.t. p_x x + p_y y = 100

\mathcal{L} = x y + \lambda (p_x x + p_y y - 100)

\frac{\partial \mathcal{L}}{\partial x} = y + \lambda p_x = 0

\frac{\partial \mathcal{L}}{\partial y} = x + \lambda p_y = 0

\frac{\partial \mathcal{L}}{\partial \lambda} = p_x x + p_y y - 100 = 0

1 = x y + \lambda (px x + py y - 100)

1x = \partial_x1 == 0

1y = \partial_y1 == 0

11 = \partial_\lambda1 == 0

soln = Solve[{1x, 1y, 11}, {x, y, \lambda}]

x y + (-100 + px x + py y) \lambda

y + px \lambda == 0

x + py \lambda == 0

-100 + px x + py y == 0

{x + y y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y + y
```

The elasticity of the functions are:

Exercise 3

Corbae 2.5.3

We write the relations in matrix form for convenience:

```
The \leq relation) on the set \{0, 1, 2, 3,4\}:
1 1 1 1 1
 0 1 1 1 1
 0 0 0 1 1
```

We can see that the matrix is upper traingular. For every element in {0, 1, 2, 3,4} either aRb or bRa. We also see that the diagonals are all 1, therefore it is reflexive.

```
The = relation on \{0, 1, 2, 3, 4\}:
 1 0 0 0 0
 0 1 0 0 0
 0 0 1 0 0
```

We can see that this relation is not complete because (for example) neither (0, 1) nor (1, 0) are in the set of =. This relationship is transitive because it is only defined for elements for identical elements.

```
The < relation on {0, 1, 2, 3, 4}:
0 1 1 1 1
 0 0 1 1 1
 0 0 0 1 1
 0 0 0 0 1
```

This relation is not complete because (0, 0) is not in the set of >.

The relation is transitive because for any aRb ∧ bRc ⇒ aRc (we can see this in the matrix representation, where all elements to the right of the first "1" in a row are also 1.

```
0 1 1 1 1 1
```

We can see the relation is not complete because elements such as (0, 0) are not in the set of ≠. We can see the set is not transitive because 3 ± 1 , 1 ± 3 does not imply 3 ± 3 (that is, (1,3) and (3, 1) are in the set of the relation \neq , but (1, 1) is not).

For example 2.5.6 we have:

	chicken	fish	beef	pork
pork	1	0	0	1
beef	0	0	1	1
fish	0	1	1	1
chicken	1	1	1	0

The following relations are in the set ≤ {(pork, chicken), (pork, pork), (beef, beef), (beef, pork), (fish, fish), (fish, beef), (fish, pork), (chicken, chicken), (chicken, fish), (chicken, beef)}. We can see that from the set {pork, beef, fish, chicken} either $a \le b$ or $b \le a$.

2.5.5

~ is rational means ~ is complete and transitive.

This means $x \ge x$ by completeness. Therefore, we can rule out $\neg [x \ge x]$. This is only possible if $x \sim x$. Suppose \sim is not transitive. This implies $\exists_{x, y, z \in \mathbb{A}} x \sim y$, $y \sim z$ but $\neg (x \sim z)$. This means $(x > z) \lor (z < z)$. Thus we have $x \sim y > z$. But we have $y \sim z$. Thus we have a contratdiction.

2.5.8

We saw previously the relation is complete. However, the relation is not transitive. For example, we see "pork R chicken", and "chicken R fish", but no "pork R fish". Thus it cannot be an equivalence relations.

2.7.6

The equivalence relation are the set of all vectors with the same length and direction.

2.7.10

$$\sim = (\gtrsim \bigcap \gtrsim^{-1}) ?$$

Exercise 4

P 72

1 a. Negate "Everyone who is majoring in math has a friend who needs help with his homework."

Let A be the set of people majoring in math

Let B be the set of people who need help with homework

Let xFy be the relation x and y are friends

We can write the above symbolically as

$$\forall_x \in A \exists_v \in B \land (x, y) \in F$$

Negating this:

 $\neg (\forall_x \in A \exists_y \in B \text{ such that } (x, y) \in F)$

 \Leftrightarrow $(\exists_x \in A \neg (\exists_y \in B \text{ such that } (x,y) \in F)$

$$\Leftrightarrow$$
 $(\exists_x \in A \text{ such that } \forall_y \in B \neg (x, y) \in F)$
 \Leftrightarrow $(\exists_x \in \text{ such that } A \forall_y \in B \neg (x, y) \notin F)$

Which we can rewrite this as "there is someone who is majoring in math who doesn't have any friends who needs help with math".

1 b. Negate "Everyone has a roommate who dislikes everyone."

Let xDy be the relation x dislikes y

Let xRy be the relation x and y are roommates

We can rewrite the above statement symbolically as:

$$\begin{aligned} &\forall_x \; \exists_y \; \text{s.t.} \; (x,y) \in \mathsf{R} \; \land \; (x,y) \in \mathsf{D} \\ &\neg (\forall_x \; \exists_y \; \text{s.t.} \; (x,y) \in \mathsf{R} \; \land \; (x,y) \in \mathsf{D}) \\ &(\exists_x \; \neg (\exists_y \; \text{s.t.} \; (x,y) \in \mathsf{R} \; \land \; (x,y) \in \mathsf{D})) \\ &(\exists_x \; \text{s.t.} \; (\forall_y \; \neg ((x,y) \in \mathsf{R} \; \land \; (x,y) \in \mathsf{D}))) \\ &(\exists_x \; \text{s.t.} \; (\forall_y \; (\neg (x,y) \in \mathsf{R} \; \lor \; \neg (x,y) \in \mathsf{D}))) \\ &(\exists_x \; \text{s.t.} \; (\forall_y \; ((x,y) \in \mathsf{R} \; \lor \; (x,y) \in \mathsf{D}))) \end{aligned}$$

There is someone who has a roommate who doesn't dislike everyone.

1 c.

Negate A
$$\cup$$
 B \subseteq C \ D \neg (A \cup B \subseteq C \ D) $\Leftrightarrow \neg$ A \cap \neg B \subseteq C \D) $\Leftrightarrow \neg$ A \cap B \nsubseteq (C\D)

1 d.

Negate
$$\exists_x \forall_y [y > x \Rightarrow \exists_z (z^2 + 5 z = y)]$$

 $\forall_x \exists_y [y > x \Rightarrow \forall_z (z^2 + 5 z \neq y)]$

2.

a. There is someone in the freshman class who doesn't have a roommate.

Everyone in the freshman class has a roommate

b. Everyone likes someone, but no one likes everyone

There is someone who doesn't like anyone or there is someone who likes everyone

c.
$$\forall_a \in A \exists_b \in B \ (a \in C \Leftrightarrow b \in C)$$

 $\exists_a \in A \ \forall_b \in B \ \neg (a \in C \Leftrightarrow b \in C)$
 $\exists_a \in A \ \forall_b \in B \ (a \in C \land b \notin C) \lor (a \notin C \land b \in C)$
d. $\forall_y > 0 \ \exists_x (a \ x^2 + b \ x + c = y)$
 $\exists_x > 0 \ \forall_x \ (a \ x^2 + b \ x + c \neq y)$

3.

a. False.

Consider the number 2

There are no a, b, c in the natural number $\mathbb{N} = \{1, 2, 3...\}$ such that $a^2 + b^2 + c^2 = 2$ (however, the statement is false if 0 is in the natural numbers).

b. False

$$E_{!x}(x-4)^2=9$$

$$(x-4)^2-9=0$$

$$x == 1 || x == 7$$

c.
$$\exists_{!x} (x-4)^2 = 25$$

False,
$$x = 1$$
 or $x = 9$

d. True. From previous example, if x = 1 and y = 9 then $(x - 4)^2 = 25$ and $(y - 4)^2 = 25$

4. Prove $\neg \forall_x P(x) \iff \exists_x \neg P(x)$

$$\neg \exists_x P(x) \Leftrightarrow \forall_x \neg P(x)$$

negate both side

$$\neg(\neg \exists_x P(x)) \Longleftrightarrow \neg (\forall_x \neg P(x))$$

distribute

$$E_x \neg P(x) \iff (\neg A_x P(x))$$

5. $\neg E_x \in A P(x) \iff \forall_x \neg P(x)$

by quantifier negation law

P. 170

$$A = \{1, 2, 3\}$$

$$B = \{1, 4\}$$

$$C = \{3, 4\}$$

$$D = \{5\}$$

B
$$\cup$$
 C = {1, 3, 4}

$$B \cap C = \{4\}$$

$$A \cup C = \{1, 2, 3, 4\}$$

$$A \cap C = \{3\}$$

$$B \cup D = \{1, 4, 5\}$$

$$B \cap D = \{\}$$

$$A \times B = \{(1, 1), (1, 4), (2, 1), (2, 4), (3, 1), (3, 4)\}$$

$$A \times C = \{(1, 3), (1, 4), (2,3), (2, 4), (3, 3), (3, 4)\}$$

$$C \times D = \{(3, 5), (4, 5)\}$$

$$\mathsf{A}\times(\mathsf{B}\cap\mathsf{C})=(\mathsf{A}\times\mathsf{B})\cap(\mathsf{A}\times\mathsf{C})$$

$$\{1, 2, 3\} \times \{4\} = \{(1, 1), (1, 4), \{2, 1), (2, 4), (3, 1), (3, 4)\} \cap \{(1, 3), (1, 4), (2, 3), (2, 4), (3, 3), (3, 4)\}$$

 $\{(1, 4), (2, 4), (3, 4)\} = \{(1, 4), (2, 4), (3, 4)\}$

$$A \times (B \cup C) = (A \times B) \cup (A \times C)$$

$$\{1, 2, 3\} \times \{1, 3, 4\} = \{(1, 1), (1, 4), \{2, 1), (2, 4), (3, 1), (3, 4)\} \cup \{(1, 3), (1, 4), (2, 3), (2, 4), (3, 3), (3, 4)\}$$

$$\{(1, 1), (1, 3), (1, 4), (2, 1), (2, 3), (2, 4), (3, 1), (3, 3), (3, 4)\} = \{(1, 1), (1, 3), (1, 4), \{2, 1), (2, 3), (2, 4), (3, 1), (3, 3), (3, 4)\}$$

$$(A \times B) \cap (C \times D) = (A \cap C) \times (B \cap D)$$

{(1, 1), (1, 4), {2, 1), (2, 4), (3, 1), (3, 4)} \cap {(3, 5), (4, 5)} = {3} \cap {\} {\}

$$(A \times B) \cup (C \times D) \subseteq (A \cup C) \times (B \cup D) \\ \{(1, 1), (1, 4), \{2, 1), (2, 4), (3, 1), (3, 4)\} \cup \{(3, 5), (4, 5)\} \subseteq \{1, 2, 3, 4\} \times \{1, 4, 5\} \\ \{(1, 1), (1, 4), \{2, 1), (2, 4), (3, 1), (3, 4), (3, 5), (4, 5)\} \subseteq \{(1, 1), (1, 4), (1, 5), (2, 1), (2, 4), (2, 4), (3, 1), (3, 4), (3, 5), (4, 1), (4, 4), (4, 5)\}$$

$$A \times \emptyset = \emptyset \times A = \emptyset$$

{1, 2, 3} × $\emptyset = \emptyset$
 $\emptyset \times \{1, 2, 3\} = \emptyset$ (see page 166).

7. If A has m elements and B has n A × B will have m × n elements.

P. 178

1 a. $\{(p, q) \in P \times P \mid \text{the person p is a parent of the person q}\}$, where P is the set of all living people} The domain is the set of all living people.

The range is the set of all living people with a living parent.

b.
$$\{(x, y) \mathbb{R}^2 \mid y > x^2\}$$

The doman is \mathbb{R} . The range is all non-negative real numbers.

2 a. $\{(p,q) \in P \times P \mid \text{the person p is a brother of the person q} \} P \text{ is the set of all living people.}$ The domain is the set of all living people. The range is the set of all living people with a living brother.

b.
$$\{(x, y) \in \mathbb{R}^2 \mid y^2 = 1 - 2/(x^2 + 1)\}.$$

The domain is the set of all real numbers. the range is -1 < y < 1.

Exercise 5

The theorem and proof are not correct. Consider an empty relation (i.e. no two elements of the nonempty set are in the relation). Then R is transmitive and symmetric, but not reflexive. in order to fix the proof, we must add we must add a quantifier condition (that is for every element a in the set, there exists at least one b such that aRb).

Exercise 6

```
if relation R is irreflexive than \forall_a \neg (aRa).
If relation R is antisymmetric, then \forall_{a,b} (aRb \land bRa) \Rightarrow a = b
Therefore \forall_{a,b} ((aRb \land bRa) \Rightarrow a = b) \land ¬(aRa) \Leftrightarrow (¬(aRb \land bRa) \lor a = b) \land ¬ (aRa) \Leftrightarrow (¬aRb \lor ¬bRa \lor
a = b) ∧ ¬ aRa⇔aRb ⇒ ¬bRa
```

Exercise 7

R is symmetric: iff $a_{ij} = 1$ then $a_{ji} = 1$ in the bolean matrix representation (where a_{ij} is the i^{th} row and j^{th} column of the boolean matrix representation of the relationship).

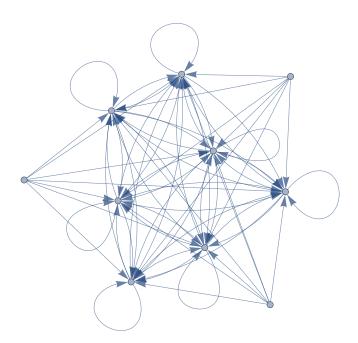
Asymettric relation: iff $a_{ij} = 1$ then $a_{ij} = 0$.

If R is complete, then the matrix is triangular.

Computational Exercise 1

```
ClearAll[result]
subsym[x_] := Module[{result = Table[Table[0, Length[x]], Length[x]]},
  Do[If[x[[i]][[j]] = 1 && x[[j]][[i]] = 1,
    result[[i]][[j]] = 1,
    0], {i, 1, Length[x]}, {j, 1, Length[x]}];
  result
 1
x = Table[{1, 0, 1, 1, 0, 1, 1, 1, 0, 1}, 10];
x // Labeled[TableForm[#], "Original Relation"] &
AdjacencyGraph[x, PlotLabel → "Graph of original adjacency matrix"]
s = subsym[x];
s // Labeled[TableForm[#], "Symmetric Subelation"] &
AdjacencyGraph[s, PlotLabel > "Symmetric Subrelation"]
1
                                            1
                   0
   0 1 1 0 1 1 0 1
0 1 1 0 1 1 0 1
0 1 1 0 1 1 1 0 1
0 1 1 0 1 1 1 0 1
0 1 1 0 1 1 1 0 1
1
1
1
    0 1 1 0
                       1
1
             1 0 1
1 0 1
                            1
    0 1
             Original Relation
```

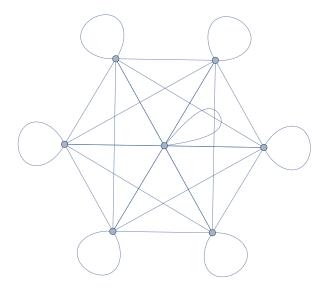
Graph of original adjacency matrix



1	0	1	1	0	1	1	1	0	-1
Τ.	U	Τ.	Ţ	0	Τ.	1	Ţ	U	Τ.
0	0	0	0	0	0	0	0	0	0
1	0	1	1	0	1	1	1	0	1
1	0	1	1	0	1	1	1	0	1
0	0	0	0	0	0	0	0	0	0
1	0	1	1	0	1	1	1	0	1
1	0	1	1	0	1	1	1	0	1
1	0	1	1	0	1	1	1	0	1
0	0	0	0	0	0	0	0	0	0
1	0	1	1	0	1	1	1	0	1

Symmetric Subelation

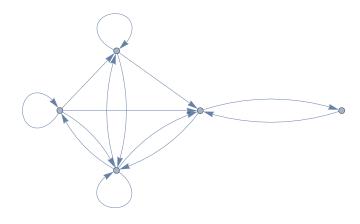
Symmetric Subrelation



Compuational Exercise 2

```
adjMatrix = Table[RandomChoice[{0, 1}, 5], 5]
adjMatrix // Labeled["Adjacency Matrix", TableForm[#]] &
adjGraph = AdjacencyGraph[adjMatrix, PlotLabel → "Adjacency Graph"]
\{\{1, 0, 1, 1, 0\}, \{1, 1, 1, 1, 0\}, \{0, 0, 0, 1, 1\}, \{1, 1, 1, 1, 0\}, \{0, 0, 1, 0, 0\}\}
 Adjacency Matrix
     0
         1
              1
1
1
          1
               1
     1
0
          0
               1
1
               1
                    0
          1
         1
               0
```

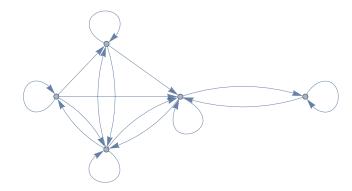
Adjacency Graph



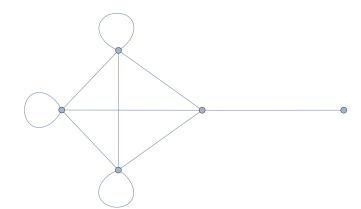
```
StringForm["Is the relation reflexive? ``",
 AllTrue[Diagonal[adjMatrix], # == 1 &]]
symmetric[x_] := Module[{},
  Do[
   Ιf
    [x[[i]][[j]] = 1 && x[[j]][[i]] = 0,
    Return[False]
   ];
   Return[True],
   {i, 1, Length[x]},
   {j, 1, Length[x]}
  ]
StringForm["Is the relation symmetric? ``", symmetric[adjMatrix]]
Is the relation reflexive? False
Is the relation symmetric? True
```

```
reflexivize[x_] := Module[{result = x},
  Do[
   result[[i]][[i]] = 1,
   {i, 1, Length[x]}];
  Return[result]]
reflexiveAdjMatrix = reflexivize[adjMatrix];
StringForm["Is the new relation (with reflexive closure) reflexive? ``",
 AllTrue[Diagonal[reflexiveAdjMatrix], # == 1 &]]
AdjacencyGraph[reflexiveAdjMatrix]
```

Is the new relation (with reflexive closure) reflexive? True

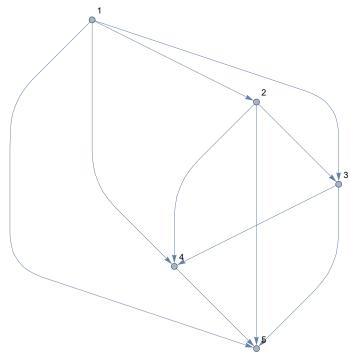


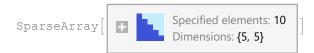
```
symmetrize[x_] := Module[{result = x}, Do[
   If[result[[i]][[j]] == 1, result[[j]][[i]] = 1,],
    \{ \texttt{i, 1, Length[result]} \}, \ \{ \texttt{j, 1, Length[result]} \} ]; 
  Return[result]]
symmetricAdjMatrix = symmetrize[adjMatrix];
StringForm["Is the new relation (with symmtric closure)? ``",
 symmetric[symmetricAdjMatrix]]
AdjacencyGraph[symmetricAdjMatrix]
Is the new relation (with symmtric closure)? True
```



Computational Exercise 3

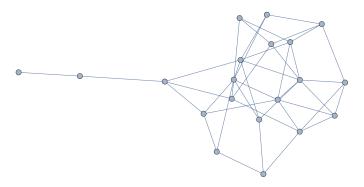
```
g = RelationGraph[Less, Range[1, 5]]
a = AdjacencyMatrix[g]
```





		<	<		
	1	2	3	4	5
1	0	1	1	1	1
2	0	0	1	1	1
2	0	0	0	1	1
4	0	0	0	0	1
5	0	0	0	0	0

g0 = RandomGraph[BernoulliGraphDistribution[20, 0.2]]



d = VertexDegree[g0]

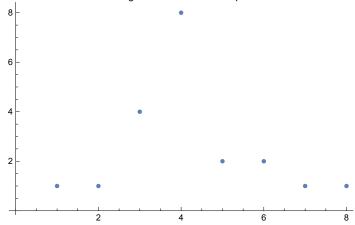
{3, 4, 5, 4, 8, 6, 7, 6, 4, 4, 4, 3, 5, 4, 4, 3, 1, 2, 4, 3}

Tally[d]

 $\{\{3, 4\}, \{4, 8\}, \{5, 2\}, \{8, 1\}, \{6, 2\}, \{7, 1\}, \{1, 1\}, \{2, 1\}\}$

ListPlot[Tally[d], PlotLabel -> "Degree Distribution of Graph"]

Degree Distribution of Graph



Computational Exercise 4

```
\texttt{adjList} \ = \ \{1 \mapsto 1 \,, \ 1 \mapsto 3 \,, \ 2 \mapsto 1 \,, \ 2 \mapsto 2 \,, \ 2 \mapsto 3 \,, \ 3 \mapsto 2 \,, \ 3 \mapsto 3 \,,
      3 \leftrightarrow 5, 4 \leftrightarrow 1, 4 \leftrightarrow 2, 4 \leftrightarrow 3, 4 \leftrightarrow 4, 5 \leftrightarrow 1, 5 \leftrightarrow 2, 5 \leftrightarrow 3, 5 \leftrightarrow 5
g1 = Graph[adjList]
AdjacencyMatrix[g1] // MatrixForm
 \{1 \leftrightarrow 1, 1 \leftrightarrow 3, 2 \leftrightarrow 1, 2 \leftrightarrow 2, 2 \leftrightarrow 3, 3 \leftrightarrow 2, 3 \leftrightarrow 3, \}
  3 \leftrightarrow 5, 4 \leftrightarrow 1, 4 \leftrightarrow 2, 4 \leftrightarrow 3, 4 \leftrightarrow 4, 5 \leftrightarrow 1, 5 \leftrightarrow 2, 5 \leftrightarrow 3, 5 \leftrightarrow 5
```

