

## Semantics for the $\lambda$ -calculus



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## CHAPTER 1

# Definitions

### 1. Algebraic Theories

DEFINITION 1 (algebraic theory). We define an algebraic theory  $T$  to be a sequence of sets  $T_n$  indexed over  $\mathbb{N}$  with for all  $1 \leq i \leq n$  elements ("variables" or "projections")  $x_{n,i} : T_n$  (we usually leave  $n$  implicit), together with a substitution operation

$$- \bullet - : T_m \times T_n^m \rightarrow T_n$$

for all  $m, n$ , such that

$$\begin{aligned} x_j \bullet g &= g_j \\ f \bullet (x_{l,i})_i &= f \\ (f \bullet g) \bullet h &= f \bullet (g_i \bullet h)_i \end{aligned}$$

for all  $1 \leq j \leq l$ ,  $f : T_l$ ,  $g : T_m^l$  and  $h : T_n^m$ .

DEFINITION 2 (algebraic theory morphism). A morphism  $F$  between algebraic theories  $T$  and  $T'$  is a sequence of functions  $F_n : T_n \rightarrow T'_n$  (we usually leave the  $n$  implicit) such that

$$\begin{aligned} F_n(x_j) &= x_j \\ F_n(f \bullet g) &= F_m(f) \bullet (F_n(g_i))_i \end{aligned}$$

for all  $1 \leq j \leq n$ ,  $f : T_m$  and  $g : T_n^m$ .

REMARK 1. We can construct binary products of algebraic theories, with sets  $(T \times T')_n = T_n \times T'_n$ , variables  $(x_i, x'_i)$  and substitution

$$(f, f') \bullet (g, g') = (f \bullet g, f' \bullet g').$$

In the same way, the category of algebraic theories has all limits.

### 2. Algebras

DEFINITION 3 (algebra). An algebra  $A$  for an algebraic theory  $T$  is a set  $A$ , together with an action

$$\bullet : T_n \times A^n \rightarrow A$$

for all  $n$ , such that

$$\begin{aligned} x_j \bullet a &= a_j \\ (f \bullet g) \bullet a &= f \bullet (g_i \bullet a)_i \end{aligned}$$

for all  $j$ ,  $f : T_m$ ,  $g : T_n^m$  and  $a : A^n$ .

DEFINITION 4 (algebra morphism). For an algebraic theory  $T$ , a morphism  $F$  between  $T$ -algebras  $A$  and  $A'$  is a function  $F : A \rightarrow A'$  such that

$$F(f \bullet a) = f \bullet (F(a_i))_i$$

for all  $f : T_n$  and  $a : A^n$ .

REMARK 2. The category of algebras has all limits. The set of a limit of algebras is the limit of the underlying algebras.

REMARK 3. Note that for an algebraic theory  $T$ , the  $T_n$  are all algebras for  $T$ , with the action given by  $\bullet$ .

### 3. Presheaves

DEFINITION 5 (presheaf). A presheaf  $P$  for an algebraic theory  $T$  is a sequence of sets  $P_n$  indexed over  $\mathbb{N}$ , together with an action

$$\bullet : P_m \times T_n^m \rightarrow P_n$$

for all  $m, n$ , such that

$$\begin{aligned} t \bullet (x_{l,i})_i &= t \\ (t \bullet f) \bullet g &= t \bullet (f_i \bullet g)_i \end{aligned}$$

for all  $t : P_l$ ,  $f : T_m^l$  and  $g : T_n^m$ .

DEFINITION 6 (presheaf morphism). For an algebraic theory  $T$ , a morphism  $F$  between  $T$ -presheaves  $P$  and  $P'$  is a sequence of functions  $F_n : P_n \rightarrow P'_n$  such that

$$F_n(t \bullet f) = F_m(t) \bullet f$$

for all  $t : P_m$  and  $f : T_n^m$ .

We will write  $PT$  for the category of  $T$ -presheaves and their morphisms.

REMARK 4. The category of presheaves has all limits. The  $n$ th set  $\bar{P}_n$  of a limit  $\bar{P}$  of presheaves  $P_i$  is the limit of the  $n$ th sets  $P_{i,n}$  of the presheaves in the limit diagram.

### 4. $\lambda$ -theories

DEFINITION 7 ( $\lambda$ -theory). A  $\lambda$ -theory is an algebraic theory  $L$ , together with sequences of functions  $\lambda_n : L_{n+1} \rightarrow L_n$  and  $\rho_n : L_n \rightarrow L_{n+1}$ , such that

$$\begin{aligned} \lambda(f) \bullet g &= \lambda(f \bullet (g_1, \dots, g_m, x_{n+1})) \\ \rho(f \bullet g) &= \rho(f) \bullet (g_1, \dots, g_m, x_{n+1}) \end{aligned}$$

for all  $f : L_m$  and  $g : L_n^m$ . ((**TODO**) : Fix)

DEFINITION 8 ( $\beta$ - and  $\eta$ -equality). We say that a  $\lambda$ -theory  $L$  satisfies  $\beta$ -equality (or that it is a  $\lambda$ -theory with  $\beta$ ) if  $\rho_n \circ \lambda_n = \text{id}_{L_n}$  for all  $n$ . We say that it satisfies  $\eta$ -equality if  $\lambda_n \circ \rho_n = \text{id}_{L_{n+1}}$  for all  $n$ .

DEFINITION 9 ( $\lambda$ -theory morphism). A morphism  $F$  between  $\lambda$ -theories  $L$  and  $L'$  is an algebraic theory morphism  $F$  such that

$$\begin{aligned} F_n(\lambda_n(f)) &= \lambda_n(F_{n+1}(f)) \\ \rho_n(F_n(g)) &= F_{n+1}(\rho_n(g)) \end{aligned}$$

for all  $f : L_{n+1}$  and  $g : L_n$ .

REMARK 5. The category of lambda theories has all limits, with the underlying algebraic theory of a limit being the limit of the underlying algebraic theories.

A  $\lambda$ -theory algebra or presheaf is a presheaf for the underlying algebraic theory.

## CHAPTER 2

### Lemmas

#### 1. The endomorphism theory

DEFINITION 10 (Endomorphism theory). Suppose that we have a category  $C$  and an object  $X : C$ , such that all powers  $X^n$  of  $X$  are also in  $C$ . The endomorphism theory  $E(X)$  of  $X$  is the algebraic theory given by  $E(X)_n = C(X^n, X)$  with projections as variables  $x_{n,i} : X^n \rightarrow X$  and a substitution that sends  $f : X^m \rightarrow X$  and  $g_1, \dots, g_m : X^n \rightarrow X$  to  $f \circ \langle g_i \rangle_i : X^n \rightarrow X^m \rightarrow X$ .

LEMMA 1.  $E(X)$  is indeed an algebraic theory.

PROOF. For  $1 \leq j \leq l$ ,  $f : E(X)_l$ ,  $g : E(X)_m^l$  and  $h : E(X)_n^m$ , we have

$$\begin{aligned} x_j \bullet g &= x_j \circ \langle g_i \rangle_i = g_j, \\ f \bullet (x_{l,i})_i &= f \circ \langle x_{l,i} \rangle_i = f \circ \text{id}_{X^l} = f \end{aligned}$$

and

$$(f \bullet g) \bullet h = f \circ \langle g_i \rangle_i \circ \langle h_i \rangle_i = f \circ \langle g_i \circ \langle h_{i'} \rangle_{i'} \rangle_i = f \bullet (g_i \bullet h)_i.$$

□

DEFINITION 11 (Endomorphism  $\lambda$ -theory). Now, suppose that the exponential object  $X^X$  exists, and that we have morphisms back and forth  $\text{abs} : X^X \rightarrow X$  and  $\text{app} : X \rightarrow X^X$ . Let, for  $Y : C$ ,  $\varphi_Y$  be the isomorphism  $C(X \times Y, X) \xrightarrow{\sim} C(Y, X^X)$ . We can give  $E(X)$  a  $\lambda$ -theory structure by setting, for  $f : E(X)_{n+1}$  and  $g : E(X)_n$ ,

$$\lambda(f) = \text{abs} \circ \varphi_{X^n}(f) \quad \rho(g) = \varphi_{X^n}^{-1}(\text{app} \circ g).$$

LEMMA 2.  $E(X)$  is indeed a  $\lambda$ -theory.

PROOF. Note that  $\varphi : C(- \times X, X) \xrightarrow{\sim} C(-, X^X)$  is a natural isomorphism, so for  $g : E(X)_n^m$ , the following diagram commutes

$$\begin{array}{ccc} C(X^m \times X, X) & \xrightarrow{- \circ (\langle g_i \rangle_i \times \text{id}_X)} & C(X^n \times X, X^X) \\ \varphi_{X^m}^{-1} \left( \begin{array}{c} \uparrow \\ \downarrow \end{array} \right) \varphi_{X^m} & & \varphi_{X^n}^{-1} \left( \begin{array}{c} \uparrow \\ \downarrow \end{array} \right) \varphi_{X^n} \\ C(X^m, X^X) & \xrightarrow{- \circ \langle g_i \rangle_i} & C(X^n, X^X) \end{array}$$

and note that  $\langle g_i \rangle_i \times \text{id}_X = \langle g_1, \dots, g_m, x_{n+1} \rangle$ . Then we have, for all  $f : E(X)_m$

$$\begin{aligned} \lambda_m(f) \bullet g &= \text{abs} \circ \varphi_{X^m}(f) \circ \langle g_i \rangle_i \\ &= \text{abs} \circ \varphi_{X^n}(f \circ \langle g_1, \dots, g_m, x_{n+1} \rangle) \\ &= \lambda_n(f \bullet (g_1, \dots, g_m, x_{n+1})) \end{aligned}$$

and

$$\begin{aligned} \rho_n(f \bullet g) &= \varphi_{X^n}^{-1}(\text{app} \circ f \circ \langle g_i \rangle_i) \\ &= \varphi_{X^m}^{-1}(\text{app} \circ f) \circ \langle g_1, \dots, g_m, x_{n+1} \rangle \\ &= \rho_m(f) \bullet (g_1, \dots, g_m, x_{n+1}). \end{aligned}$$

□

## 2. The theory presheaf

DEFINITION 12 (The theory presheaf). Let  $T$  be an algebraic theory. We can turn  $T$  into an  $T$ -presheaf  $\tilde{T}$  by setting  $\tilde{T}_n = T_n$  and using the substitution from  $T$ :

$$\bullet : \tilde{T}_m \times T_n^m \rightarrow \tilde{T}_n.$$

LEMMA 3.  $\tilde{T}$  is indeed a presheaf.

PROOF. For all  $t : \tilde{T}_l$ ,  $f : T_m^l$  and  $g : T_n^m$ ,

$$t \bullet (x_{l,i})_i = t$$

and

$$(t \bullet f) \bullet g = t \bullet (f_i \bullet g)_i$$

because  $T$  is an algebraic theory.  $\square$

LEMMA 4. Given an algebraic theory  $T$  and a  $T$ -presheaf  $Q$ , we have for all  $n$  a bijection of sets

$$\varphi : PT(\tilde{T}^n, Q) \cong Q_n.$$

PROOF. Take  $\varphi(f) = f_n(x_1, \dots, x_n)$ .

Conversely, take  $\varphi^{-1}(q)$  to be the presheaf morphism that sends  $t : T_m^n$  to  $q \bullet t : Q_m$ . This is indeed a presheaf morphism, since for all  $t : T_l^n$  and  $f : T_m^l$ ,

$$\varphi^{-1}(q)(t \bullet f) = q \bullet t \bullet f = \varphi^{-1}(q)(t) \bullet f.$$

Now, for a presheaf morphism  $f : T^n \rightarrow Q$  and  $t : T_m^n$ , we have

$$\varphi^{-1}(\varphi(f))(t) = f_n(x_1, \dots, x_n) \bullet t = f_n((x_1, \dots, x_n) \bullet t) = f_n(t_1, \dots, t_n) = f_n(t).$$

Conversely, given  $q : Q_n$ , we have

$$\varphi(\varphi^{-1}(q)) = q \bullet (x_1, \dots, x_n) = q.$$

which concludes the proof.  $\square$

## 3. The ‘+l’ presheaf

Let  $\iota_{m,n} : T_m \rightarrow T_{m+n}$  denote the function that sends  $f$  to  $f \bullet (x_{m+n,1}, \dots, x_{m+n,m})$ . Note that

$$\iota_{m,n}(f) \bullet g = f \bullet (g_i)_{i \leq m}$$

and

$$\iota_{m,n}(f \bullet g) = f \bullet g \bullet (x_i)_i = f \bullet (g_i \bullet (x_j)_j)_i = f \bullet (\iota_{m,n}(g_i))_i.$$

For tuples  $x : X^m$  and  $y : X^n$ , let  $x+y$  denote the tuple  $(x_1, \dots, x_m, y_1, \dots, y_n) : X^{m+n}$ .

DEFINITION 13 (The ‘+l’ presheaf). Given a  $T$ -presheaf  $Q$ , we can construct a presheaf  $A(Q, l)$ , given by  $A(Q, l)_n = Q_{n+l}$ . Then, for  $q : A(Q, l)_m$  and  $f : T_n^m$ , the substitution is given by

$$q \bullet_{A(Q,l)} f = q \bullet_Q ((\iota_{n,l}(f_i))_i + (x_{n+i})_i)$$

LEMMA 5. The  $+l$  presheaf is a presheaf

PROOF. We have, for  $q : A(Q, l)_n$ ,

$$\begin{aligned} q \bullet_{A(Q,l)} (x_i)_i &= q \bullet_Q ((\iota_{n,l}(x_i))_i + (x_{n+i})_i) \\ &= q \bullet_Q ((x_i)_i + (x_{n+i})_i) \\ &= q \bullet_Q (x_i)_i \\ &= q. \end{aligned}$$



We have, for  $q : A(Q, k)_l$ ,  $f : T_m^l$  and  $g : T_n^m$ ,

$$\begin{aligned}
q \bullet_{A(Q, k)} f \bullet_{A(Q, k)} g &= q \bullet_Q ((\iota_{m, l}(f_i))_i + (x_{m+i})_i) \bullet_Q ((\iota_{n, l}(g_i))_i + (x_{n+i})_i) \\
&= q \bullet_Q (((\iota_{m, l}(f_i) \bullet_T ((\iota_{n, l}(g_j))_j + (x_{n+j})_j))_i + (x_{m+i} \bullet_T ((\iota_{n, l}(g_j))_j + (x_{n+j})_j))_i)) \\
&= q \bullet_Q ((f_i \bullet_T (\iota_{n, l}(g_j))_j)_i + (x_{n+i})_i) \\
&= q \bullet_Q ((\iota_{n, l}(f_i \bullet_T g))_i + (x_{n+i})_i) \\
&= q \bullet_{A(Q, k)} (f_i \bullet_T g).
\end{aligned}$$

□

#### 4. Exponentiability of the theory presheaf

LEMMA 6. *For all  $l$ , the presheaf  $\tilde{T}^l$  is exponentiable.*

PROOF. We will show that  $A(-, l)$  constitutes a right adjoint to the functor  $- \times \tilde{T}^l$ . We will do this using universal arrows ([ML98], Chapter IV.1, Theorem 2 (iv)). To that end, we will need for all  $Q : PT$  a universal arrow  $\varphi : A(Q, l) \times \tilde{T}^l \rightarrow Q$ .

For  $q : A(Q, l)_n = Q_{n+l}$  and  $t : \tilde{T}_n^l$ , we take  $\varphi(q, t) = q \bullet_Q ((x_{n,i})_i + t)$ .

This is a presheaf morphism, since for all  $q : A(Q, l)_m^l$ ,  $t : \tilde{T}_m^l$  and  $f : T_n^m$ ,

$$\begin{aligned}
\varphi((q, t) \bullet_{A(Q, l) \times \tilde{T}^l} f) &= \varphi(q \bullet_{A(Q, l)} f, t \bullet_{\tilde{T}^l} f) \\
&= q \bullet_{A(Q, l)} f \bullet_Q ((x_i)_i + (t \bullet_{\tilde{T}^l} f)) \\
&= q \bullet_Q ((\iota_{n, l}(f_i))_i + (x_{n+i})_i) \bullet_Q ((x_i)_i + (t \bullet_{\tilde{T}^l} f)) \\
&= q \bullet_Q ((\iota_{n, l}(f_i) \bullet_T ((x_j)_j + (t \bullet_{\tilde{T}^l} f)))_i + (x_{n+i} \bullet_T ((x_j)_j + (t \bullet_{\tilde{T}^l} f))))_i \\
&= q \bullet_Q ((f_i \bullet_T (x_j)_j)_i + ((t \bullet_{\tilde{T}^l} f)_i)_i) \\
&= q \bullet_Q ((f_i)_i + (t_i \bullet_{\tilde{T}^l} f)_i) \\
&= q \bullet_Q ((x_i \bullet_T f)_i + (t_i \bullet_T f)_i) \\
&= q \bullet_Q ((x_i)_i + t) \bullet_Q f \\
&= \varphi(q, t) \bullet_Q f.
\end{aligned}$$

Now, given any presheaf  $Q' : PT$  we need to show that any morphism  $\psi : Q' \times \tilde{T}^l \rightarrow Q$  factors uniquely as  $\varphi \circ (\tilde{\psi} \times \text{id}_{\tilde{T}^l})$  for some  $\tilde{\psi} : Q' \rightarrow A(Q, l)$ .

So, given such a  $\psi$ , and given  $q : Q'_n$ , we take  $\tilde{\psi}(q) = \psi(\iota_{n, l}(q), (x_{n+i})_i)$

This is a presheaf morphism, since for all  $q : Q'_m$  and  $f : T_n^m$ ,

$$\begin{aligned}
\tilde{\psi}(q \bullet f) &= \psi(\iota_{n, l}(q \bullet f), (x_{n+i})_i) \\
&= \psi(q \bullet (\iota_{n, l}(f_i))_i, (x_{n+i})_i) \\
&= \psi((\iota_{m, l}(q), (x_{m+i})_i) \bullet_{Q' \times \tilde{T}^l} ((\iota_{n, l}(f_i))_i + (x_{n+i})_i)) \\
&= \psi(\iota_{m, l}(q), (x_{m+i})_i) \bullet_Q ((\iota_{n, l}(f_i))_i + (x_{n+i})_i) \\
&= \tilde{\psi}(q) \bullet_{A(Q, l)} f.
\end{aligned}$$

Note that indeed  $\varphi \circ (\tilde{\psi} \times \text{id}_{\tilde{T}^l}) = \psi$ :

$$\begin{aligned}
\varphi(\tilde{\psi}(q), t) &= \varphi(\psi(\iota_{n, l}(q), (x_{n+i})_i), t) \\
&= \psi(\iota_{n, l}(q), (x_{n+i})_i) \bullet_Q ((x_i)_i + t) \\
&= \psi(\iota_{n, l}(q) \bullet_Q ((x_i)_i + t), (x_{n+i})_i \bullet_Q ((x_i)_i + t)) \\
&= \psi(q \bullet (x_i)_i, (t_i)_i) \\
&= \psi(q, t).
\end{aligned}$$

Now, suppose that we have another  $\tilde{\psi}' : Q' \rightarrow A(Q, l)$  such that  $\varphi \circ (\tilde{\psi}' \times \text{id}_{\tilde{T}^l}) = \psi$ . Then we have

$$\begin{aligned}
 \tilde{\psi}(q) &= \psi(\iota_{n,l}(q), (x_{n+i})_i) \\
 &= (\varphi \circ (\tilde{\psi}' \times \text{id}_{\tilde{T}^l}))(\iota_{n,l}(q), (x_{n+i})_i) \\
 &= \varphi(\tilde{\psi}'(\iota_{n,l}(q)), (x_{n+i})_i) \\
 &= \tilde{\psi}'(\iota_{n,l}(q)) \bullet ((x_i)_i + (x_{n+i})_i) \\
 &= \iota_{n,l}(\tilde{\psi}'(q)) \bullet ((x_i)_i + (x_{n+i})_i) \\
 &= \tilde{\psi}'(q) \bullet (x_i)_i \\
 &= \tilde{\psi}'(q),
 \end{aligned}$$

so  $\tilde{\psi}$  is unique, which completes the proof.  $\square$

Now, this adjunction  $- \times \tilde{T}^l \dashv A(-, l)$  induces a natural isomorphism

$$\varphi : PT(- \times \tilde{T}^l, \tilde{T}) \xrightarrow{\sim} PT(-, A(\tilde{T}, l))$$

LEMMA 7. For all  $f : PT(\tilde{T}^n \times \tilde{T}^l, \tilde{T})$ ,

$$\varphi_{\tilde{T}^n}(f)(q) = f(\iota_{m,l}(q), (x_{m+i})_i)$$

PROOF. **(TODO)**  $\square$

LEMMA 8. For all  $f : PT(\tilde{T}^n, A(\tilde{T}, l))$ ,

$$\varphi_{\tilde{T}^n}^{-1}(f)(q, t) = f(q) \bullet ((x_i)_i + t).$$

PROOF. **(TODO)**  $\square$

## CHAPTER 3

# Theorems

### 1. Scott's Representation Theorem

**THEOREM 1.** *Any  $\lambda$ -theory  $L$  is isomorphic to the endomorphism  $\lambda$ -theory  $E(\tilde{L})$  of  $\tilde{L}$  in the presheaf category of  $L$ .*

**PROOF.** First of all, remember that  $\tilde{L}$  is indeed exponentiable and that  $\tilde{L}^{\tilde{L}} = A(\tilde{L}, 1)$ . Now, since  $L$  is a  $\lambda$ -theory, we have functions back and forth  $\lambda : A(\tilde{L}, 1) \rightarrow \tilde{L}$  and  $\rho : \tilde{L} \rightarrow A(\tilde{L}, 1)$ . These are presheaf morphisms because for all  $f : A(\tilde{L}, 1)_m$  and  $g : \tilde{L}_m$  and  $t : T_n^m$ ,

$$\lambda(f \bullet_{A(\tilde{L}, 1)} t) = \lambda(f \bullet_{\tilde{L}} ((\iota_{m,1}(t_i))_i + (x_{n+1}))) = \lambda(f) \bullet_{\tilde{L}} t$$

and

$$\rho(g \bullet_{\tilde{L}} t) = \rho(g) \bullet_{\tilde{L}} ((\iota_{m,1}(t_i))_i + (x_{n+1})) = \rho(g) \bullet_{A(\tilde{L}, 1)} t.$$

Therefore,  $E(\tilde{L})$  is indeed a  $\lambda$ -theory.

For any presheaf  $Q$  and for any  $n$ , we have a bijection  $PL(L^n, Q) \cong Q_n$ . Then we have  $\varphi : E(\tilde{L})_n \cong L_n$ . This bijection is an isomorphism of  $\lambda$ -theories, since it preserves the  $x_i$ ,  $\bullet$ ,  $\rho$  and  $\lambda$ : for all  $1 \leq j \leq n$ ,  $f : E(\tilde{L})_m$ ,  $g : E(\tilde{L})_{m+1}$  and  $h : E(\tilde{L})_n^m$ .

$$\begin{aligned} \varphi(x_j) &= x_j(x_1, \dots, x_n) \\ &= x_j; \\ \varphi(f \bullet h) &= f \circ \langle h_i \rangle_i((x_i)_i) \\ &= f((h_i((x_j)_j))_i) \\ &= f((x_i)_i \bullet (h_i((x_j)_j))_i) \\ &= f((x_i)_i) \bullet (h_i((x_j)_j))_i \\ &= \varphi(f) \bullet (\varphi(h_i))_i; \\ \varphi(\rho(f)) &= \rho(f)((x_i)_i) \\ &= \rho(f((x_i)_i)) \bullet (x_i)_i \\ &= \rho(f((x_i)_i)) \\ &= \rho(\varphi(f)); \\ \varphi(\lambda(g)) &= \lambda(g)((x_i)_i) \\ &= \lambda(\varphi_{X^n}(g)((x_i)_i)) \\ &= \lambda(g(\iota_{m,l}((x_i)_i) + (x_{m+1}))) \\ &= \lambda(g((x_i)_i)) \\ &= \lambda(\varphi(g)). \end{aligned}$$

□

## 2. Locally cartesian closedness of the category of retracts

DEFINITION 14 (Category of retracts). The category of retracts for a  $\lambda$ -theory  $L$  is the category with objects  $f : L_n$  such that  $f \bullet f = f$  and it has as morphisms  $g : f \rightarrow f'$  the terms  $g : L_n$  such that  $f' \bullet g \bullet f = g$ . The object  $f : L_n$  has identity element  $f$ , and we have composition  $g \circ g' = g \bullet g'$ . These are morphisms (**TODO**)

LEMMA 9. *The category of retracts is indeed a category.*

PROOF. (**TODO**) □

THEOREM 2. *The category of retracts is locally cartesian closed (TODO) .*

## 3. The Fundamental Theorem of the $\lambda$ -calculus

DEFINITION 15 ( $\Lambda$ ). There is a special  $\lambda$ -theory, given by the  $\lambda$ -calculus itself.  $\Lambda_n$  is the set of  $\lambda$ -terms with  $n$  free variables, the  $x_i$  are the free variables, and  $\bullet$  is given by substitution.  $\lambda$  sends  $f : \Lambda_{n+1}$  to  $\lambda x_{n+1}. f$  and  $\rho$  sends  $f : \Lambda_n$  to  $\iota_{n,1}(f)x_{n+1}$  in  $\Lambda_n$ .

LEMMA 10.  *$\Lambda$  is indeed a  $\lambda$ -theory.*

PROOF. (**TODO**) □

LEMMA 11.  *$\Lambda$  is the initial  $\lambda$ -theory.*

PROOF. Given a  $\lambda$ -theory  $L$ , we construct a morphism  $f : \Lambda \rightarrow L$  by induction on the  $\lambda$ -terms. We set  $f(x_i) = x_i$ ,  $f(\lambda(t)) = \lambda(f(t))$  and  $f(st) = \rho(f(s)) \bullet ((x_i)_i + (f(t)))$ .

This is a  $\lambda$ -theory morphism because (**TODO**)

It is unique, since (**TODO**) □

DEFINITION 16 (Pullback of algebras). If we have a morphism of algebraic theories  $f : T' \rightarrow T$ , we have a functor  $AT \rightarrow AT'$ .

On objects, it sends a  $T$ -algebra  $A$  to a  $T'$ -algebra with set  $A$  and action  $g \bullet_{T'} a = f(g) \bullet_T a$ . This is a  $T'$ -algebra because (**TODO**) .

On morphisms, it sends  $\varphi : A \rightarrow A$  to  $\varphi : A \rightarrow A$ . This is a  $T'$ -algebra morphism because for all  $g : T'_n$  and  $a : A^n$ , we have

$$\varphi(g \bullet_{T'} a) = \varphi(f(g) \bullet_T a) = f(g) \bullet_T \varphi(a) = g \bullet_{T'} \varphi(a).$$

LEMMA 12. *This is indeed a functor.*

PROOF. (**TODO**) □

DEFINITION 17 (Term algebra). Given an algebraic theory  $T$ , for every  $n$ ,  $T_n$  together with the action operator  $\bullet : T_m \times T_n^m \rightarrow T_n$  gives a  $T$ -algebra.

LEMMA 13.  *$T_n$  is indeed a  $T$ -algebra.*

PROOF. (**TODO**) □

DEFINITION 18. For all  $n$ , we have a functor from lambda theories to  $\Lambda$ -algebras. It sends the  $\lambda$ -theory  $L$  to the  $L$ -algebra  $L_n$  and then turns this into a  $\Lambda$ -algebra via the morphism  $\Lambda \rightarrow L$ .

It sends morphisms  $f : L \rightarrow L'$  to  $f_n : L_n \rightarrow L'_n$ . This is a  $\Lambda$ -algebra morphism because (**TODO**)

LEMMA 14. *This indeed constitutes a functor.*

PROOF. (**TODO**) □

REMARK 6. Note that for a monoid  $M$ , if we view  $M$  as a category, the category  $[M^{\text{op}}, \mathbf{SET}]$  consists of sets with a right  $M$ -action.

DEFINITION 19 (The exponential object in the presheaf category). Given a monoid  $M$ , if we have two presheaves (sets with right  $M$ -actions)  $P$  and  $P'$ , we have a set of  $M$ -equivariant maps

$$F_{P,P'} = \left\{ f : M \times P \rightarrow P' \mid \prod_{p:P, m,m':M} f(m,p)m' = f(mm',pm') \right\}$$

with a right  $M$ -action, given by  $f m'(m,p) = f(m'm,p)$ . This is again  $M$ -equivariant because

$$f m'(m,p)m'' = f(m'm,p)m'' = f(m'mm'',pm'') = f m'(mm'',pm''),$$

so  $F_{P,P'}$  is a presheaf.

Now, to show that  $F_{P,P'}$  is the exponential object  $P'^P$ , we show that for any  $P$ ,  $F_{P,-}$  is the left adjoint of  $- \times P$ . So we need for all  $P' : PT$ , a universal arrow  $\varphi : F_{P,P'} \times P \rightarrow P'$ .

First of all, we have an evaluation map  $\varphi : F_{P,P'} \times P \rightarrow P'$  given by  $(f,p) \mapsto f(I,p)$  for  $I$  the unit of the monoid. This map is equivariant because for all  $m$ ,

$$(f,p)m = (fm,pm) \mapsto fm(I,pm) = f(m,pm) = f(I,p)m.$$

Now, given any presheaf  $Q$  and any morphism  $\psi : Q \times P \rightarrow P'$ , take  $\tilde{\psi} : Q \rightarrow F_{P,P'}$  given by  $\tilde{\psi}(q)(m,p) = \psi(qm,p)$ . This is equivariant because

$$\tilde{\psi}(q)m(m',p) = \tilde{\psi}(q)(mm',p) = \psi(qmm',p) = \tilde{\psi}(qm)(m',p)$$

and we have

$$\varphi(\tilde{\psi}(q),p) = \tilde{\psi}(q)(I,p) = \psi(q,p).$$

Now, suppose that we have  $\tilde{\psi}' : Q \rightarrow F_{P,P'}$  such that  $\varphi \circ (\tilde{\psi}' \times \text{id}_P) = \psi$ . Then for all  $q : Q$ ,  $m : M$  and  $p : P$ ,

$$\tilde{\psi}(q)(m,p) = \psi(qm,p) = \varphi(\tilde{\psi}'(qm),p) = \tilde{\psi}'(qm)(I,p) = \psi'(q)m(I,p) = \psi'(q)(m,p),$$

so  $\tilde{\psi}$  is unique and  $F_{P,P'}$  is an exponential object.

DEFINITION 20 (n-functional terms). Let  $A$  be a  $\Lambda$ -algebra. We define

$$A(n) = \{a : A \mid (\lambda x_2 x_3 \dots x_{n+1}, x_1 x_2 x_3 \dots x_{n+1}) \bullet a = a\}.$$

DEFINITION 21. Take  $\mathbf{1}_n = (\lambda x_1 \dots x_n, x_1 \dots x_n) \bullet () : A$ .

DEFINITION 22. We define composition as  $a \circ b = (\lambda x_3, x_1(x_2 x_3)) \circ (a, b)$  for  $a, b : A$ .

LEMMA 15. *This composition is associative.*

PROOF. (TODO) □

DEFINITION 23 (The monoid of a  $\Lambda$ -algebra). Now we make  $A(1)$  into a monoid with unit  $\lambda x_1, x_1$ .

LEMMA 16. *This is indeed a monoid.*

PROOF. (TODO) □

From here on, we will assume that  $\Lambda$  (and therefore, any  $\lambda$ -theory) satisfies  $\beta$ -equality.

LEMMA 17. *For  $a : A$ ,  $a$  is in  $A(n)$  iff  $\mathbf{1}_n \circ a = a$ .*

PROOF.

$$\begin{aligned}
\mathbf{1}_n \circ a &= (\lambda x_3, x_1(x_2x_3)) \bullet (((\lambda x_1 \dots x_n, x_1 \dots x_n) \bullet ()), a) \\
&= (\lambda x_3, x_1(x_2x_3)) \bullet (((\lambda x_2 \dots x_{n+1}, x_2 \dots x_{n+1}) \bullet a), x_1 \bullet a) \\
&= ((\lambda x_3, x_1(x_2x_3)) \bullet ((\lambda x_2 \dots x_{n+1}, x_2 \dots x_{n+1}), x_1)) \bullet a \\
&= (\lambda x_2, (\lambda x_3 \dots x_{n+2}, x_3 \dots x_{n+2})(x_1x_2)) \bullet a \\
&= (\lambda x_2x_3 \dots x_{n+1}, x_1x_2 \dots x_{n+1}) \bullet a.
\end{aligned}$$

□

DEFINITION 24 (The presheaf category of a  $\Lambda$ -algebra). Let  $A$  be a  $\Lambda$ -algebra. If we view the monoid  $A(1)$  as a one-object category, we define the category  $PA$  to be the category of presheaves  $[A(1)^{\text{op}}, \mathbf{SET}]$ .

DEFINITION 25 (The objects  $A(n)$  in  $PA$ ). Given  $a : A(n)$  and  $b : A(1)$ , we have

$$\mathbf{1}_n \circ (a \circ b) = (\mathbf{1}_n \circ a) \circ b = a \circ b,$$

so  $a \circ b : A(n)$  and we have a right  $A(1)$ -action on  $A(n)$ , which makes  $A(n)$  into an object in  $PA$ .

LEMMA 18. We have  $A(1)^{A(1)} \cong A(2)$ .

PROOF. We have a bijection  $\varphi : A(2) \cong F_{A(1), A(1)}$ , given by

$$\varphi(a)(b, b') = (\lambda x_4, x_1(x_2x_4)(x_3x_4)) \bullet (a, b, b').$$

Note that  $\varphi(d)$  is equivariant because **(TODO)** Now,  $\varphi$  is a presheaf morphism because **(TODO)**

Take  $p = \lambda x_1, x_1(\lambda x_2x_3, x_2)$  and  $q = \lambda x_1, x_1(\lambda x_2x_3, x_3)$ . These are elements of  $A(1)$ . Note that for terms  $c_1, c_2$

$$\begin{aligned}
p(\lambda x_1, x_1c_1c_2) &= (\lambda x_1, x_1c_1c_2)(\lambda x_2x_3, x_2) \\
&= (\lambda x_1x_3, x_2)c_1c_2 \\
&= c_1.
\end{aligned}$$

In the same way,  $q \circ (\lambda x_1x_2, x_2c_1c_2) = c_2$ .

An inverse is given by

$$\psi : f \mapsto \lambda x_1x_2, f(p, q)(\lambda x_3, x_3x_1x_2).$$

This is a presheaf morphism because **(TODO)**

This is an inverse, because given  $f : F_{A(1), A(1)}$  and  $(a_1, a_2) : A(1) \times A(1)$ , we have

$$\begin{aligned}
\varphi(\psi(f))(a_1, a_2) &= u(\lambda x_1x_2, f(p, q)(\lambda x_3, x_3x_1x_2))(a_1, a_2) \\
&= \lambda x_1, (\lambda x_2x_3, f(p, q)(\lambda x_4, x_4x_2x_3))(a_1x_1)(a_2x_1) \\
&= \lambda x_1, f(p, q)(\lambda x_2, x_2(a_1x_1)(a_2x_1)) \\
&= f(p, q) \circ (\lambda x_1, (\lambda x_2, x_2(a_1x_1)(a_2x_1))) \\
&= f(p \circ (\lambda x_1, (\lambda x_2, x_2(a_1x_1)(a_2x_1))), q \circ (\lambda x_1, (\lambda x_2, x_2(a_1x_1)(a_2x_1)))) \\
&= f(\lambda x_1, p(\lambda x_2, x_2(a_1x_1)(a_2x_1)), \lambda x_1, q(\lambda x_2, x_2(a_1x_1)(a_2x_1))) \\
&= f(\lambda x_1, a_1x_1, \lambda x_1, a_2x_1) \\
&= f(a_1, a_2).
\end{aligned}$$

The last line is because  $a_i : A(1)$  and therefore  $\lambda x_1, a_ix_1 = a_i$ .

On the other hand, if we have  $a_1 : A(2)$ , we have

$$\begin{aligned}\psi(\varphi(a_1)) &= \psi((a_2, a_3) \mapsto \lambda x_1, a_1(a_2 x_1)(a_3 x_1)) \\ &= \lambda x_1 x_2, (\lambda x_3, a_1(p x_3)(q x_3))(\lambda x_3, x_3 x_1 x_2) \\ &= \lambda x_1 x_2, a_1(p(\lambda x_3, x_3 x_1 x_2))(q(\lambda x_3, x_3 x_1 x_2)) \\ &= \lambda x_1 x_2, a_1 x_1 x_2 \\ &= a_1.\end{aligned}$$

The last line is because  $a_1 : A(2)$  and therefore  $\lambda x_1 x_2, a_1 x_1 x_2 = a_1$ .

Therefore, this map is a bijection and an isomorphism.  $\square$

**DEFINITION 26** (Endomorphism  $\lambda$ -theory of a  $\Lambda$ -algebra).  $PA$  borrows products from **SET**. Therefore, the algebraic theory  $E(A(1))$  exists. Now note that  $A(1)$  is exponentiable and  $A(1)^{A(1)} \cong A(2)$ . Note that  $A(2) \subseteq A(1)$  and that  $(\lambda x_2 x_3, x_1 x_2 x_3) \bullet -$  gives a function from  $A(1)$  to  $A(2)$ . This gives  $E(A(1))$  a  $\lambda$ -theory structure.

**DEFINITION 27** (Pullback functor on presheaves for a  $\Lambda$ -algebra). Let  $f : A \rightarrow A'$  be a  $\Lambda$ -algebra morphism. Then for all  $a : A(n)$ ,

$$\mathbf{1}_n \circ f(a) = f(\mathbf{1}_n) \circ f(a) = f(\mathbf{1}_n \circ a),$$

so we have an induced morphism  $f : A(n) \rightarrow A'(n)$ .

Now, given a presheaf  $P : PA'$ . We can create a presheaf  $f^*P : PA$  by taking the set of  $P$ , and, for  $p : P$  and  $a : A$ , setting  $pa = p \circ f(a)$ . This is indeed a presheaf because **(TODO)**

Now, given a morphism  $g : P \rightarrow P'$ , we get a morphism by taking the function on the sets of  $P$  and  $P'$ . This is a morphism because **(TODO)**

**LEMMA 19.** *The above indeed constitutes a functor.*

**PROOF.** **(TODO)**  $\square$

Left Kan extension then gives a left adjoint  $f_* : PA \rightarrow PA'$ .

**LEMMA 20.** *We have  $f_*(A(1)) \cong A'(1)$ .*

**PROOF.** **(TODO)**  $\square$

**LEMMA 21.**  *$f_*$  preserves finite products.*

**PROOF.** **(TODO)**  $\square$

**DEFINITION 28.** Since  $f_*$  preserves finite products, given an element of  $g : E(A(1))(n) = PA(A(1)^n, A(1))$ , we get

$$\#f_*(g) : PA'(f(A(1)^n), f(A(1))) \cong PA'(A'(1)^n, A'(1)) = E(A'(1))(n).$$

**LEMMA 22.**  *$\#f_* : E(A(1)) \rightarrow E(A'(1))$  is a map of  $\lambda$ -theories.*

**PROOF.** **(TODO)**  $\square$

**DEFINITION 29.** We have an isomorphism  $E(A(1))(0) \cong A$  given by  $a \mapsto aI$ .

**LEMMA 23.** *This is indeed an isomorphism of  $\Lambda$ -algebras.*

**PROOF.** **(TODO)**  $\square$

**LEMMA 24.** *Given  $g : A \rightarrow A'$ ,*

**THEOREM 3.** *There exists an adjoint equivalence between the category of  $\lambda$ -theories, and the category of algebras of  $\Lambda$ .*

PROOF. We will show that the functor  $L \mapsto L_0$  is an equivalence of categories.

It is essentially surjective, because  $L$  is isomorphic **(TODO)** to  $E(A(1))$ .

Now, given morphisms  $f, f' : L \rightarrow L'$ . Suppose that  $f_0 = f'_0$ . Suppose that  $L$  and  $L'$  have  $\beta$ -equality. Then, given  $l : L_n$ , we have

$$f_n(l) = \rho^n(\lambda^n(f_n(l))) = \rho^n(f_0(\lambda^n(l))) = \rho^n(f'_0(\lambda^n(l))) = \rho^n(\lambda^n(f'_n(l))) = f'_n(l),$$

so the functor is faithful.

The functor is full because a  $\Lambda$ -algebra morphism  $f : A \rightarrow A'$  induces a functor  $f^* : PA' \rightarrow PA$ , and via left Kan extension we get a left adjoint  $f_* : PA \rightarrow PA'$  with  $f_*(A(1)) \cong A'(1)$ . Now,  $f_*$  preserves (finite) products, so we have maps  $PA(A(1)^n, A(1)) \rightarrow PA'(A'(1)^n, A'(1))$  and so a map  $E(A(1)) \rightarrow E(A'(1))$ . This map, when restricted to a map  $PA(1, A(1)) \rightarrow PA'(1, A(1))$ , and transported along the isomorphism  $a \mapsto aI$  **(TODO)**, is equal to  $f$  **(TODO)**.  $\square$

LEMMA 25. *The category of  $T$ -algebras has coproducts.*

PROOF. **(TODO)**  $\square$

DEFINITION 30 (Theory of extensions). Let  $T$  be an algebraic theory and  $A$  a  $T$ -algebra. We can define an algebraic theory  $T_A$  called ‘the theory of extensions of  $A$ ’ with  $(T_A)_n = T_n + A$ . The left injection of the variables  $x_i : T_n$  gives the variables. Now, take  $h : (T_n + A)^m$ . Sending  $g : T_m$  to  $\varphi(g) := g \bullet h$  gives a  $T$ -algebra morphism  $T_m \rightarrow T_n + A$  since

$$\varphi(f \bullet g) = f \bullet g \bullet h = f \bullet (g_i \bullet h) = f \bullet (\varphi(g_i))_i.$$

This, together with the injection morphism of  $A$  into  $T_n + A$ , gives us a  $T$ -algebra morphism from the coproduct:  $T_m + A \rightarrow T_n + A$ . We especially have a function on sets  $(T_m + A) \times (T_n + A)^m \rightarrow T_n + A$ , which we will define our substitution to be.

LEMMA 26.  *$T_A$  is indeed an algebraic theory.*

PROOF. **(TODO)**  $\square$



## Bibliography

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