

# Partial Super Permutation Algorithm\*

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## 1 Keywords

- *array*: The array containing the super permutation
- *n*: The number we are finding partial super permutation for
- *Position*: The position of *n* in the last *n* number of term in array. So, if we are finding the super permutation for 5, and our current matrix is (1), the ***position*** will be 2.

$$\begin{pmatrix} 1, & 2, & 3, & 4, & 5 \\ 1, & 2, & 3, & 4, & 1 \\ 5, & 2, & 3, & 4, & 1 \\ 2, & 5, & 3, & 4, & 1 \end{pmatrix} \quad (1)$$

- *Shift*: The interchanging of position of *n* with what is on it's right. In (1), the position of 5(on the second last row) is shifted with 2 to get the last row.
- *Ledge Jump*: When *n* jumps from its current position to right. In (1), 5 ledge jumps on the first row.

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\*I am writing this after quite a long time, so there might be some errors. I would like to apologize.

## 2 Algorithm

Here is the explanation of the algorithm in not so good English. The first explanation will be in words, then the algorithm will be shown and then a demo. I think super permutation for 5 will be a good fit as it is short and contains most of the concepts. Again, I don't know if this algorithm has already been found, and if it has been found, I will be more than happy to remove it ☺.

### 2.1 In Words

Let  $A$  be that will be containing the superpermutation for  $n^{th}$  number. For super permutation of  $n$ , we start with a sorted array starting from 1 and ending with  $n$ . Initially,

$$A = (1, 2, 3, 4, \dots, n-1, n) \quad (2)$$

We shift the  $n$  by 1 units and check if the number so formed is in the array. So here  $n$  with shift with 1 and check if the array  $(n, 2, 3, 4, \dots, n-1, 1)$  belongs in  $A$ , if it belongs in  $A$  we shift  $n$  by another 1 unit and check if that resulting array is in  $A$  or not. Here, the second matrix doesn't belong in the matrix original  $A$  matrix so the resulting matrix will look something like this.

$$A = \begin{pmatrix} 1, & 2, & 3, & 4, & \dots, & n-1, & n \\ 1, & 2, & 3, & 4, & \dots, & n-1, & 1 \\ n, & 2, & 3, & 4, & \dots, & n-1, & 1 \end{pmatrix} \quad (3)$$

The array that is added to  $A$  is added after the position of  $n$  and the number is with which a unique array was formed is interchanged. Here, 1 and  $n$ 's interchanging created a new array, so we added the array to the matrix. What about the second array on the second?! Here, we are changing the position of  $n$  and 1 keeping all the digits the same, since 1 is after the ledge jump, the replacement should be after the jump. So,  $n$  jumps to the next row keeping 1 on first column of second row safe.

Now the same pattern holds true for the entirety of the execution,  $n$  shifts by 1 unit to the right, we check if it occurs in the array, if it doesn't we add the resulting array to our  $A$ ; however, if the resulting array is in  $A$ , we shift again to the right, we do this until a new array is found. If no array is found the execution is stopped. What if a new array is found after multiple shifting? Well in that case the new array is formed by interchanging the position of  $n$  and the digit and reversing everything in between  $n$  and the position of the number and adding it to  $A$ . So let's say in (3), no new array is found until  $n$  replaces with 4, in that case the array that is added to  $A$  will be:

$$(4, 3, 2, n, \dots, n-1, 1) \quad (4)$$

Thus the final  $A$  matrix will look something like this:

$$A = \begin{pmatrix} 1, & 2, & 3, & 4, & \cdots, & n-1, & n \\ 1, & 2, & 3, & 4, & \cdots, & n-1, & 1 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ n, & 2, & 3, & 4, & \cdots, & n-1, & 1 \\ 4, & 3, & 2, & n, & \cdots, & n-1, & 1 \end{pmatrix}$$

This same rule holds during the ledge jump as well.

## 2.2 In Algorithm

In term of algorithm it can be shown as follows:

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**Algorithm:** Partial Super Permutation

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A = (1, 2, 3, 4, ..., n-1, n);
C = ();
while All possible shift isn't in A do
    C = A[::5];
    Shift the position of n by 1 step to the right ;
    if C is not in A then
        Reverse every number between the original position and the
        unique position and add it to the original array;
    end
end
end

```

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## 2.3 C++ Code

The code can be found here [arogyad](#).

## 2.4 Demo

Lets do the demo for  $n = 5$ . We begin with:

$$(1, 2, 3, 4, 5)$$

We shift from the right hand side, since the interchanging of 1 and 5 results in a new array; we can add this new array to our older one. So the new formed array can be given as:

$$\begin{pmatrix} 1, & 2, & 3, & 4, & 5 \\ 1, & 2, & 3, & 4, & 1 \\ 5, & 2, & 3, & 4, & 1 \end{pmatrix}$$

We, then, interchange the position of 2 and 5. Since the resulting array doesn't appear in the original array before, it belongs to the original array.

$$\begin{pmatrix} 1, & 2, & 3, & 4, & 5 \\ 1, & 2, & 3, & 4, & 1 \\ 5, & 2, & 3, & 4, & 1 \\ 2, & 5, & 3, & 4, & 1 \end{pmatrix}$$

Repeating the same steps, we can form the next array as well.

$$\begin{pmatrix} 1, & 2, & 3, & 4, & 5 \\ 1, & 2, & 3, & 4, & 1 \\ 5, & 2, & 3, & 4, & 1 \\ 2, & 5, & 3, & 4, & 1 \\ 2, & 3, & 5, & 4, & 1 \end{pmatrix}$$

On to the next step, when we interchange 5 and 4, the resulting array already exists i.e (2, 3, 4, 5, 1) already exists in the array (Five digits starting from first row second column forms (2, 3, 4, 5, 1)). So, we interchange 5 and 1, which doesn't exist in the array, so the next 5 digits are (2, 3, 1, 4, 5) So the main array becomes:

$$\begin{pmatrix} 1, & 2, & 3, & 4, & 5 \\ 1, & 2, & 3, & 4, & 1 \\ 5, & 2, & 3, & 4, & 1 \\ 2, & 5, & 3, & 4, & 1 \\ 2, & 3, & 5, & 4, & 1 \\ 2, & 3, & 1, & 4, & 5 \end{pmatrix}$$

Following this algorithm, we can create the complete super permutation for  $n \leq 6$  and partial super permutation for  $n \geq 7$ .