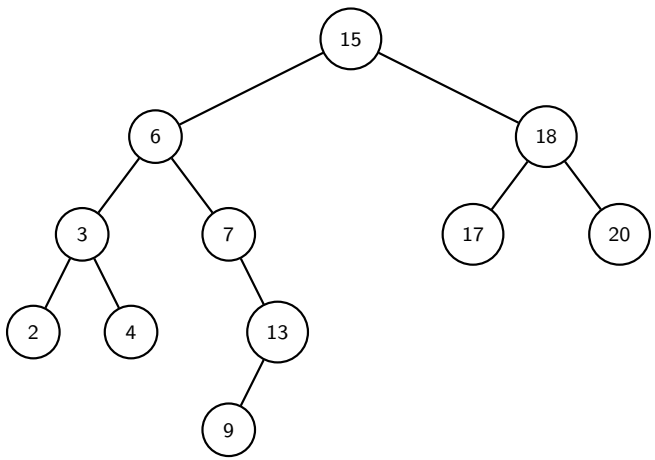
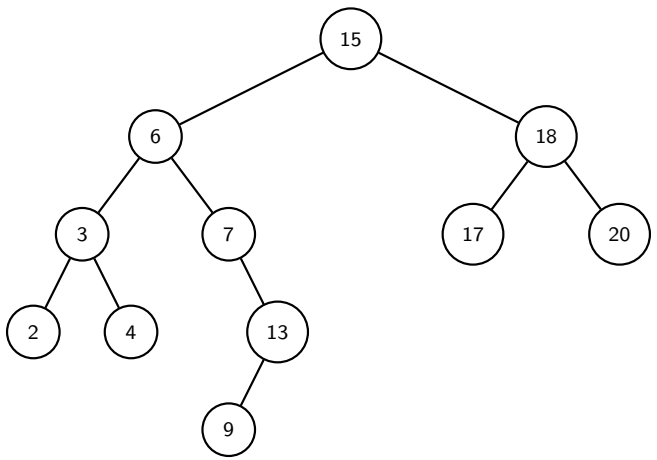
Successor

- The structure of BST allows us to determine the successor of a node without ever comparing keys!



Successor

- The structure of BST allows us to determine the successor of a node without ever comparing keys!
- Two cases may arise:
  - **Case 1:** The right sub-tree of a node x is non-empty.
    - **Successor of x:** The leftmost node in the right sub-tree.
    - $\therefore$ call TREE-MINIMUM($right[x]$).



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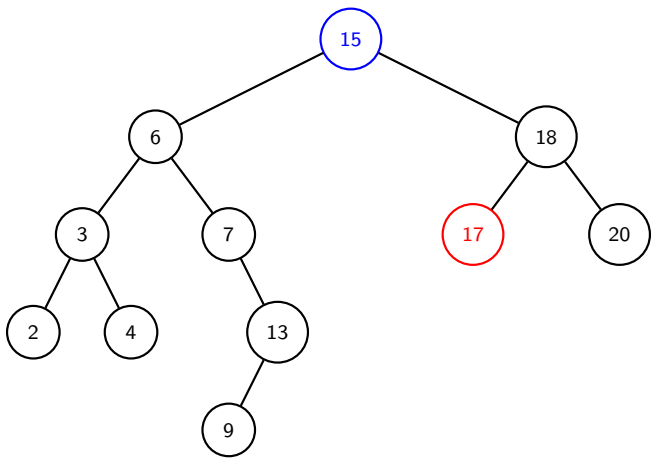

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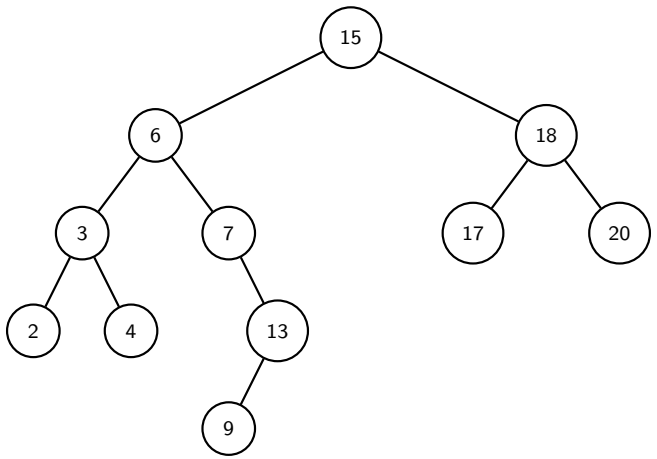

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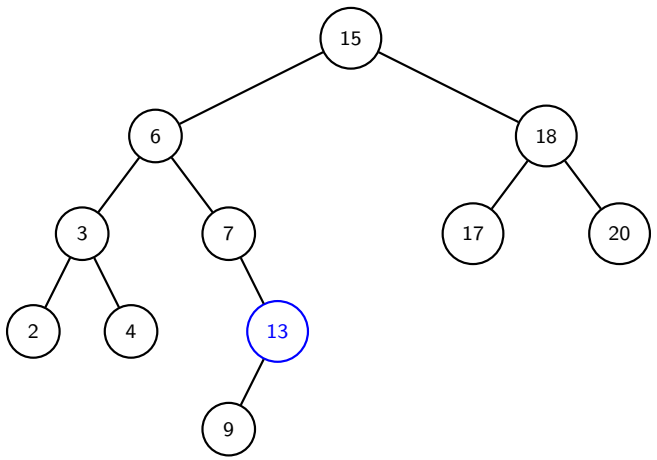
    - **Example:** Consider the node 13.



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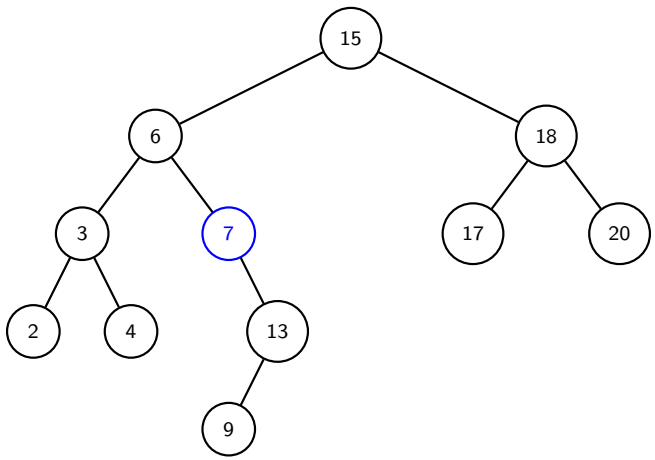

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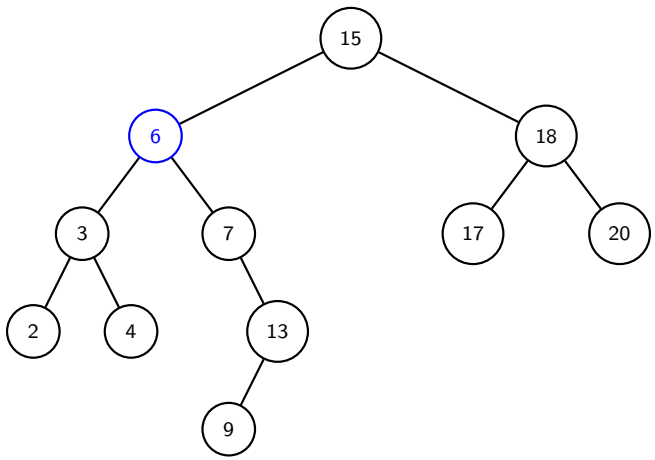

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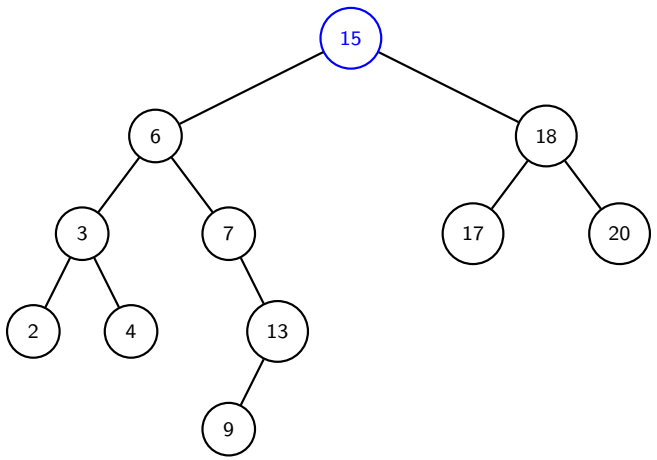

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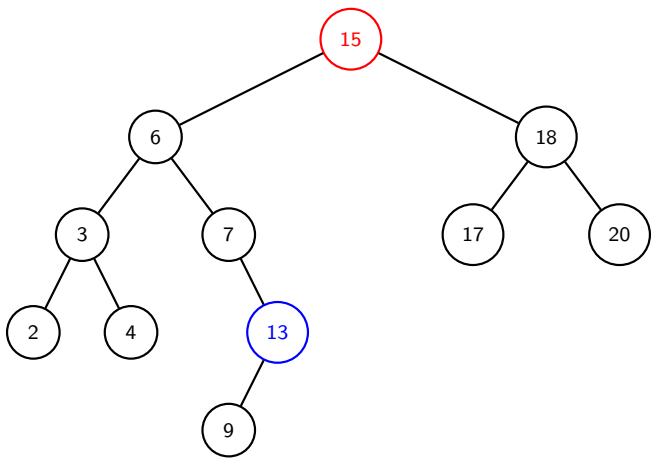
    - Then this parent is the successor.
    - **Example:** The successor of 13 is 15.



Tree-Successor

Tree-Successor(x)

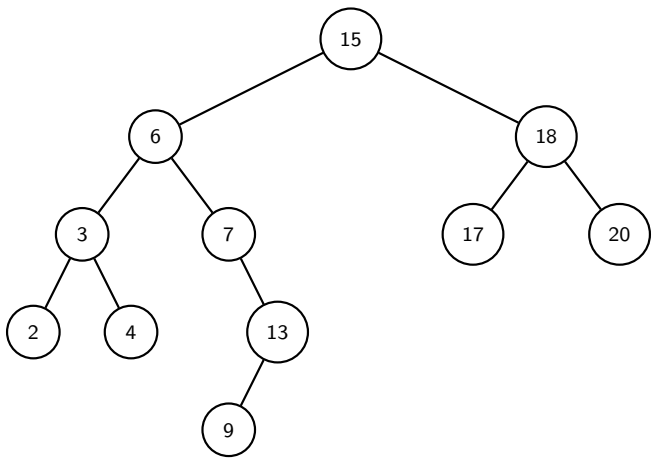
I/P: A node x whose successor we need to find.
O/P: The successor of x.

```
Begin
  if (right[x] ≠ nil)
    return Tree-Minimum(right[x]);

  y ← parent[x];
  while (y ≠ nil) and (x = right[y])
    x ← y;
    y ← parent[y];
  return y;
End
```

Predecessor

- The structure of BST allows us to determine the Predecessor of a node without ever comparing keys!



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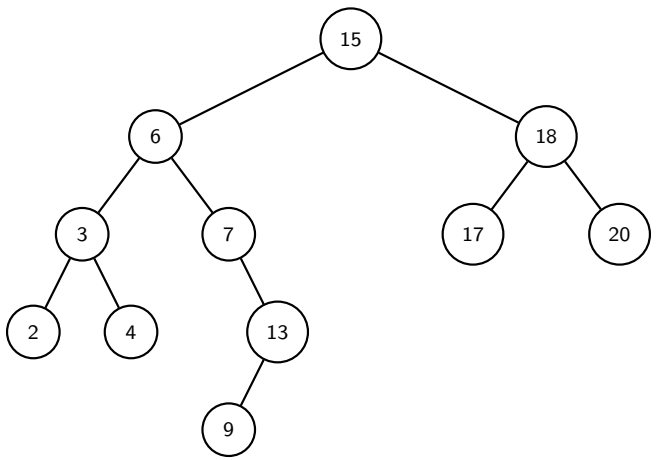
- The structure of BST allows us to determine the Predecessor of a node without ever comparing keys!
- Two cases may arise:
  - **Case 1:** The left sub-tree of a node x is non-empty.
    - **Predecessor of x:** The rightmost node in the left sub-tree.
    - $\therefore$ call TREE-MAXIMUM($left[x]$).



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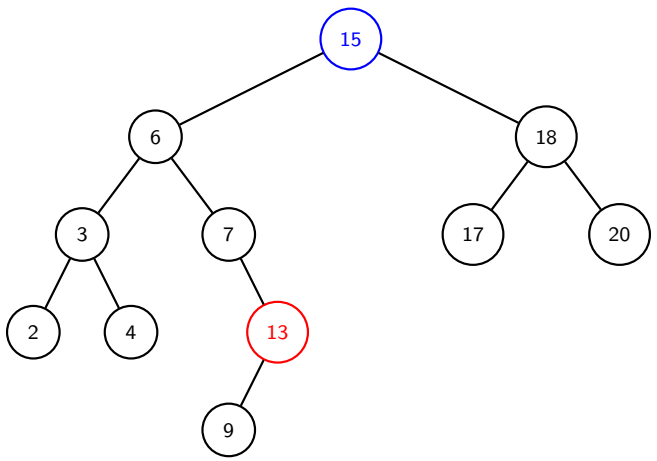
    - **Example:** Predecessor of 15 is 13.



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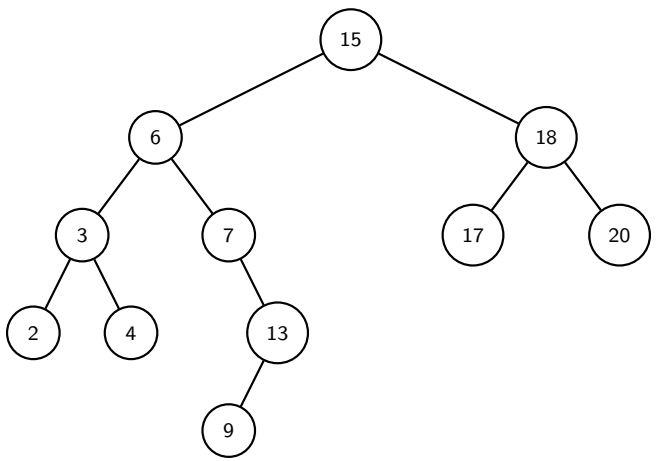

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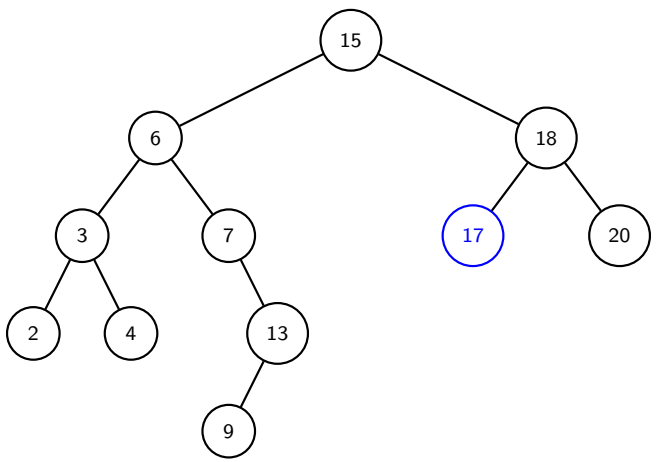

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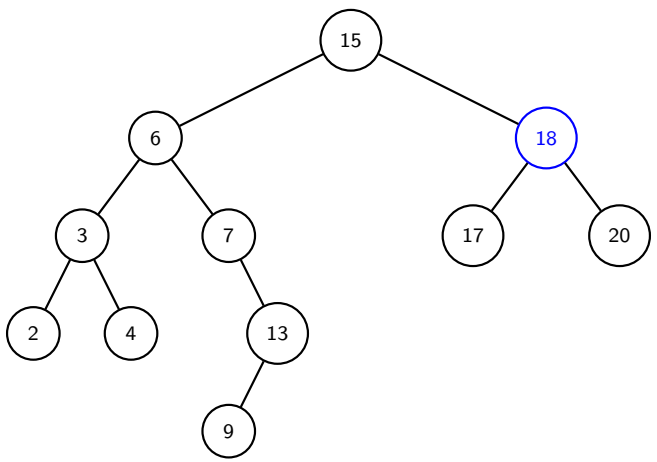

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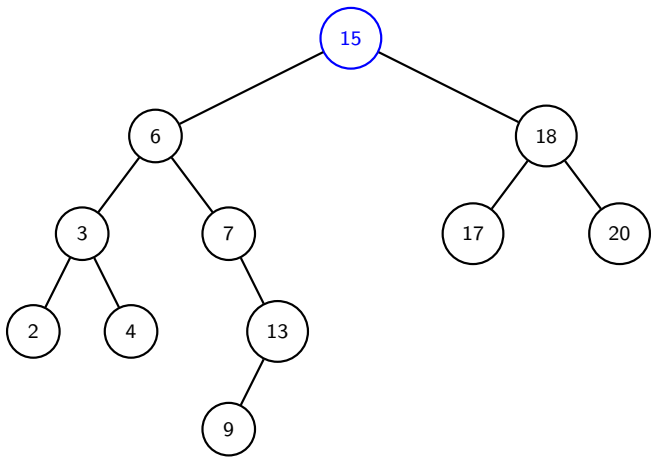

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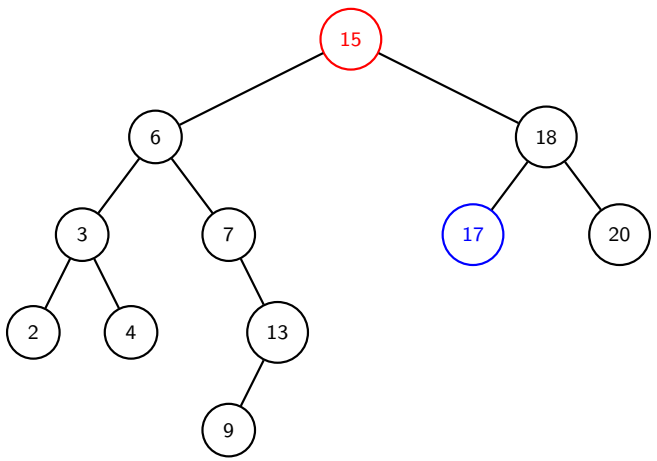

Tree-Predecessor

Tree-Predecessor(x)

I/P: A node x whose predecessor we need to find.
O/P: The predecessor of x.

```
Begin
  if (left[x] ≠ nil)
    return Tree-Maximum(left[x]);

  y ← parent[x];
  while (y ≠ nil) and (x = left[y])
    x ← y;
    y ← parent[y];
  return y;
End
```