



جامعة طرابلس - كلية تقنية المعلومات



Design and Analysis Algorithms

تصميم وتحليل خوارزميات

ITGS301

المحاضرة الثالثة: Lecture 3


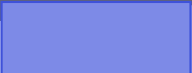





The Problem Of Sorting

Sorting

- Sorting is one of the most common data processing applications , the through which data are arranged according to their values.
- If data were not ordered , we would spend hours trying to find a single piece of information.

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- Data may be sorted in either ascending sequence or descending sequence . and if the order of the sort is not specified its assumed to be ascending..

Insertion Sort

- In the insertion sort , the list is divided in two parts : sorted and unsorted.
- In each pass the first element of the unsorted sub list is transferred to the sorted sub list by inserting it at appropriate place.
- If we have a list of n elements , it will take at most $n-1$ passes to sort the data.

```
InsertionSort(A, n) {  
    for i = 2 to n {  
        key = A[i]  
        j = i - 1;  
        while (j > 0) and (A[j] > key) {  
            A[j+1] = A[j]  
            j = j - 1  
        }  
        A[j+1] = key  
    }  
}
```

InsertionSort(A, n)

{

For $i = 2$ to n {

$Key = A[i]$

$j = i - 1$

 While ($j > 0$) and ($A[j] > key$)

 {

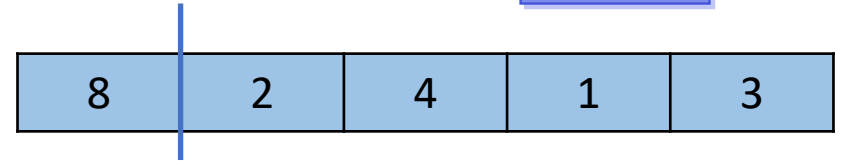
$A[j+1] = A[j]$

$j = j - 1$

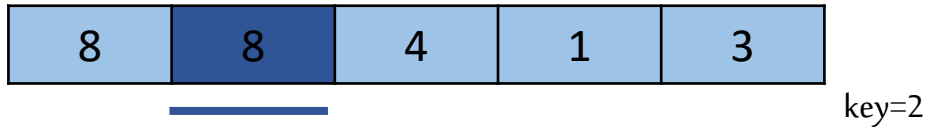
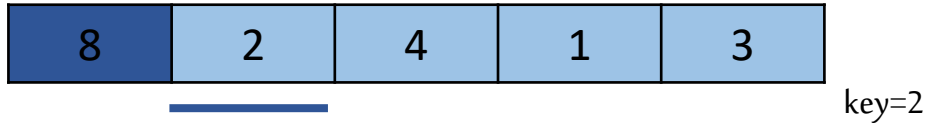
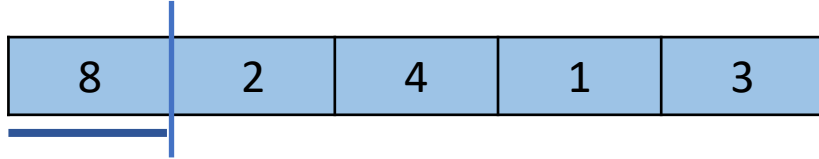
 }

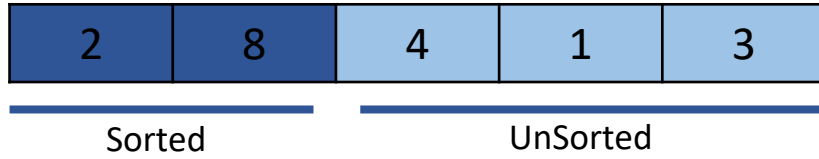
$A[j+1] = key$

}



Example 1 :





key=1

2	4	8	8	3
---	---	---	---	---

key=1

2	4	4	8	3
---	---	---	---	---

2	2	4	8	3
---	---	---	---	---

key=1

1	2	4	8	3
---	---	---	---	---

key=3



1	2	4	8	8
---	---	---	---	---

1	2	4	4	8
---	---	---	---	---

1	2	3	4	8
---	---	---	---	---

Sorted

Time Complexity

For i=2 to n {	n
Key= A [i]	$n-1$
j=i-1	$n-1$
While (j>0 & A [j] > key) {	$\sum_{i=2}^n t_i$
A [j+1] = A [j]	$\sum_{i=2}^n (t_i - 1)$
j=j-1	$\sum_{i=2}^n (t_i - 1)$
}	
A [j+1] = key	$n-1$
}	

where t is the number of while tests

Time Complexity

Statement	Time	Best case	Worst case
For i = 2 to n {	n	n	n
Key = A[i]	n-1	n-1	n-1
j= i - 1	n-1	n-1	n-1
While (j > 0) and (A[j] > key)	$\sum_{i=2}^n (ti)$	n-1	$n(n+1)/2$
A[j+1] = A[j]	$\sum_{i=2}^n (ti - 1)$	$\sum_{i=2}^n (1 - 1) = 0$	$n(n-1)/2$
j = j - 1	$\sum_{i=2}^n (ti - 1)$	$\sum_{i=2}^n (1 - 1) = 0$	$n(n-1)/2$
A[j+1] = key	n-1	n-1	n-1

Best case running time:

$$\begin{aligned} T(n) &= n + (n-1) + (n-1) + (n-1) + 0 + 0 + (n-1) \\ &= 5n - 4 \end{aligned}$$

$$T(n) = O(n)$$

Worst case running time:

$$T(n) = n + (n-1) + (n-1) + n(n+1)/2 + n(n-1)/n + n(n-1)/2 + (n-1)$$

$$T(n) = O(n^2)$$

Example 2:

8	2	4	9	3	6
---	---	---	---	---	---

8	8	4	9	3	6
---	---	---	---	---	---

2	8	4	9	3	6
---	---	---	---	---	---

2	8	8	9	3	6
---	---	---	---	---	---



2	4	8	9	3	6
---	---	---	---	---	---

2	4	8	9	3	6
---	---	---	---	---	---

2	4	8	9	9	6
---	---	---	---	---	---

2	4	8	8	9	6
---	---	---	---	---	---

2	4	4	8	9	6
---	---	---	---	---	---



2	3	4	8	9	6
---	---	---	---	---	---

2	3	4	8	9	9
---	---	---	---	---	---

2	3	4	8	8	9
---	---	---	---	---	---

2	3	4	6	8	9
---	---	---	---	---	---

The End . 