

Probably Right, Probably on Time: An Analysis of CAN in the Presence of Host and Network Faults



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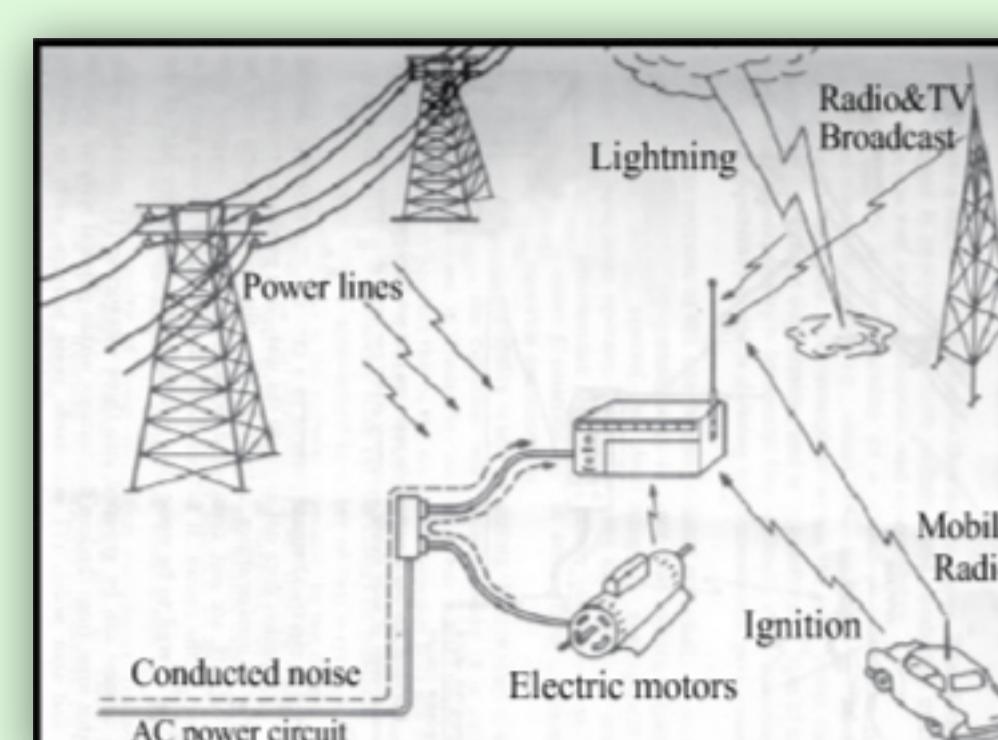


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Safety-critical RT systems susceptible to **electromagnetic interference (EMI)**

Examples

- **Automotive systems** surrounded by electric motors
- **Industrial systems** close to high-power machinery
- **Robots** operating under hard radiation

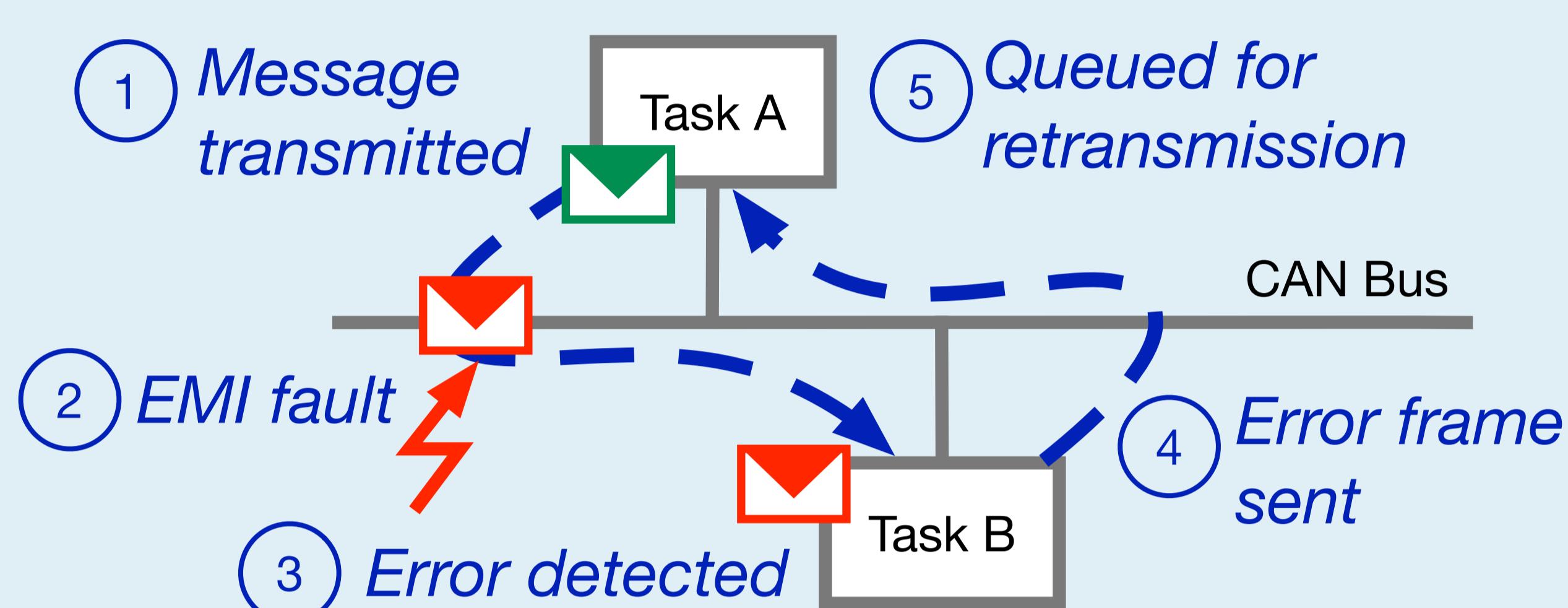


Source: http://www.bojal.com.tw/en/teach_EMI.html

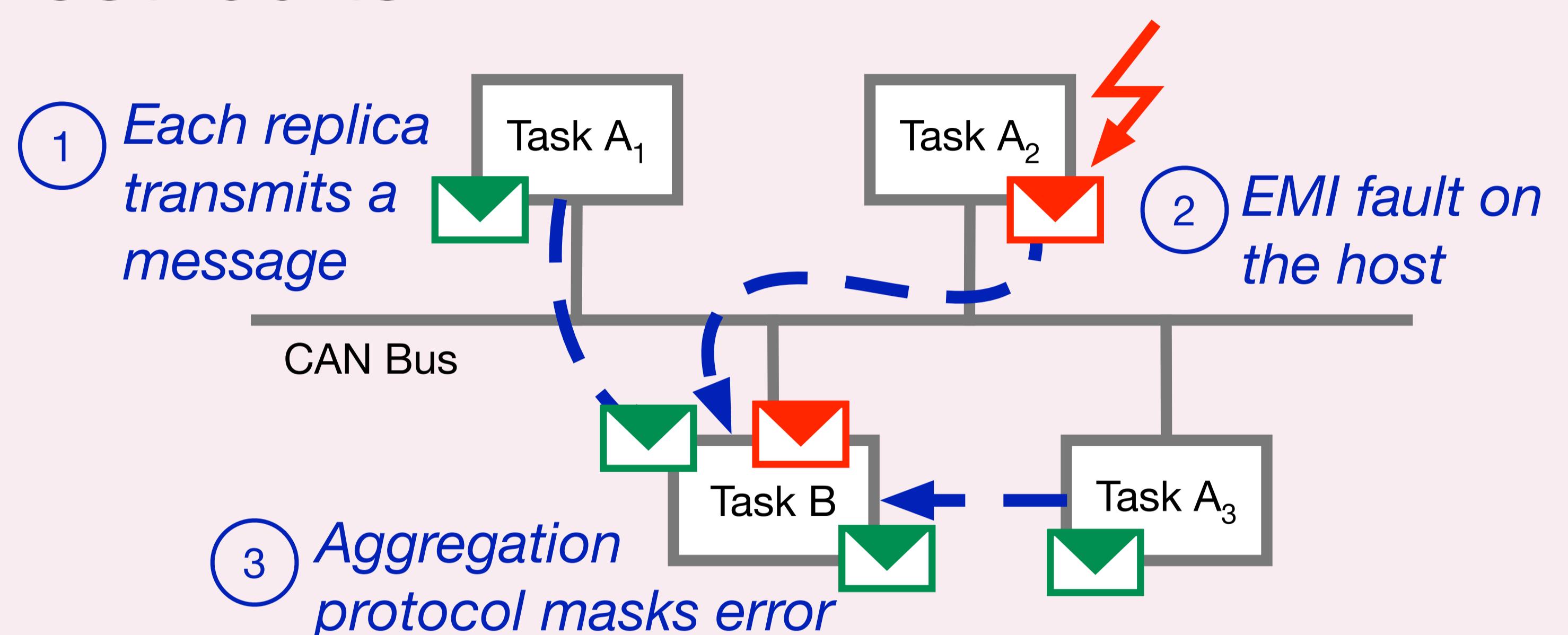
Effects of EMI

- **Host faults:** hangs, crashes, incorrect outputs
- **Transmission faults:** corrupted messages in networked systems

Retransmissions to tolerate transmission faults



Active replication of tasks to tolerate host faults



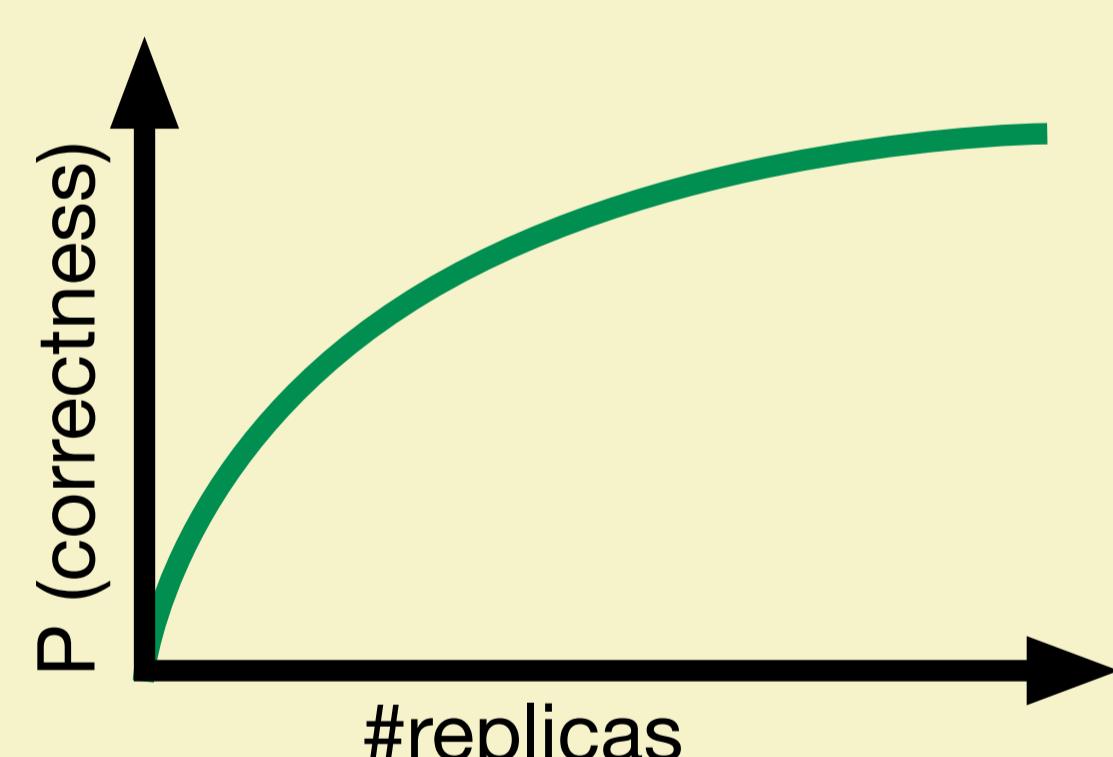
How to quantify the **inherent tradeoff** between retransmission and replication?

Higher replication

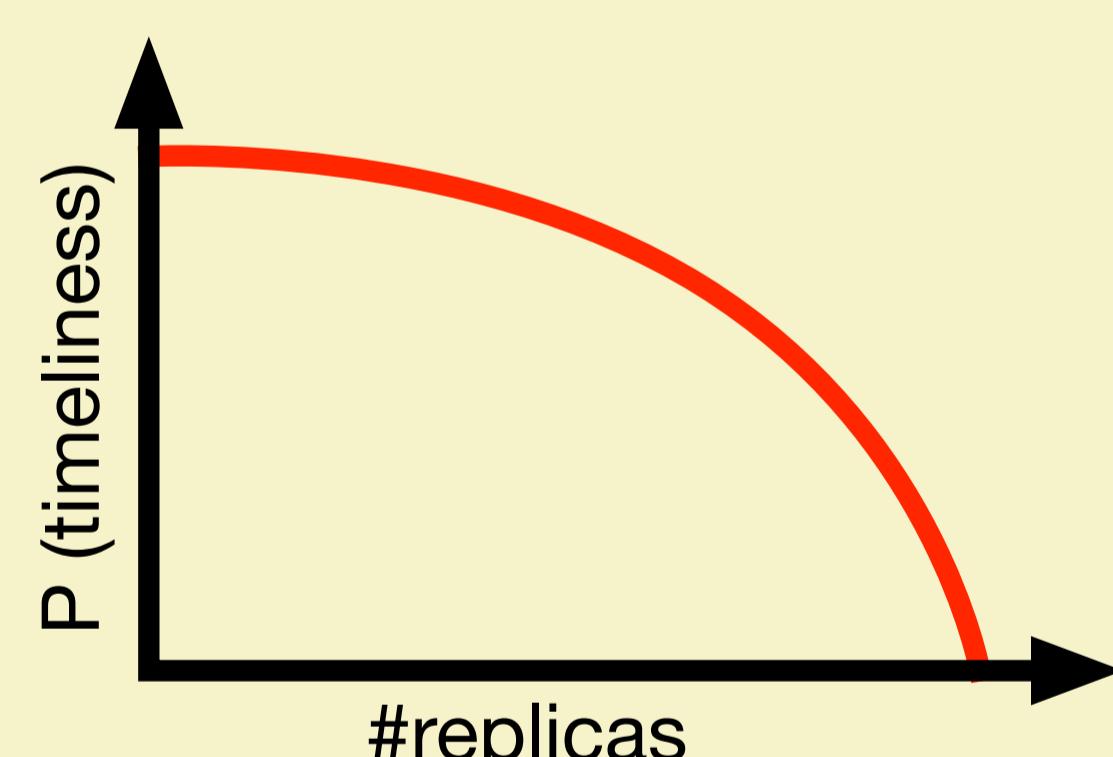
- ⇒ Better resiliency against host faults
- ⇒ Higher probability of correctness
- ⇒ But increased bus load

Increased bus load

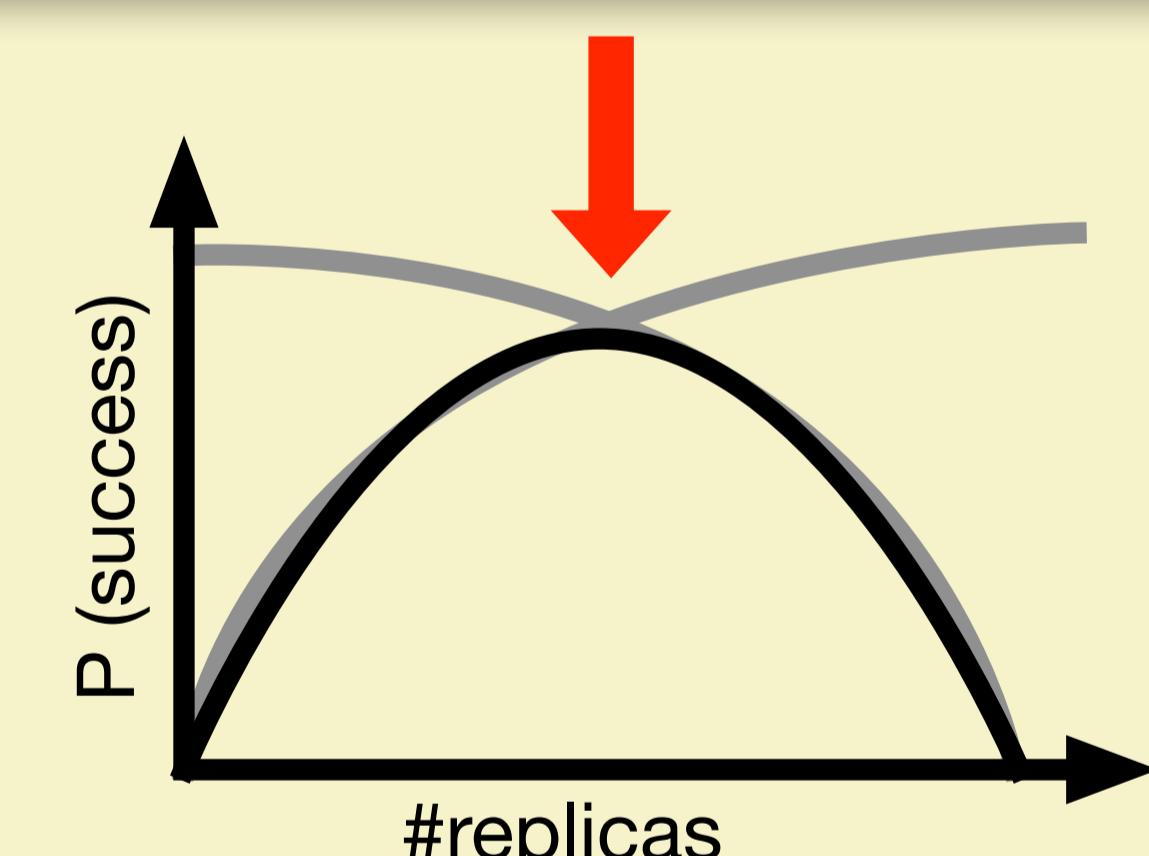
- ⇒ Less slack for retransmissions
- ⇒ More deadline violations
- ⇒ Lower probability of timeliness



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Problem
 What is the **probability of a “successful” message transmission**, i.e., that task B **infers the correct message value on time**?

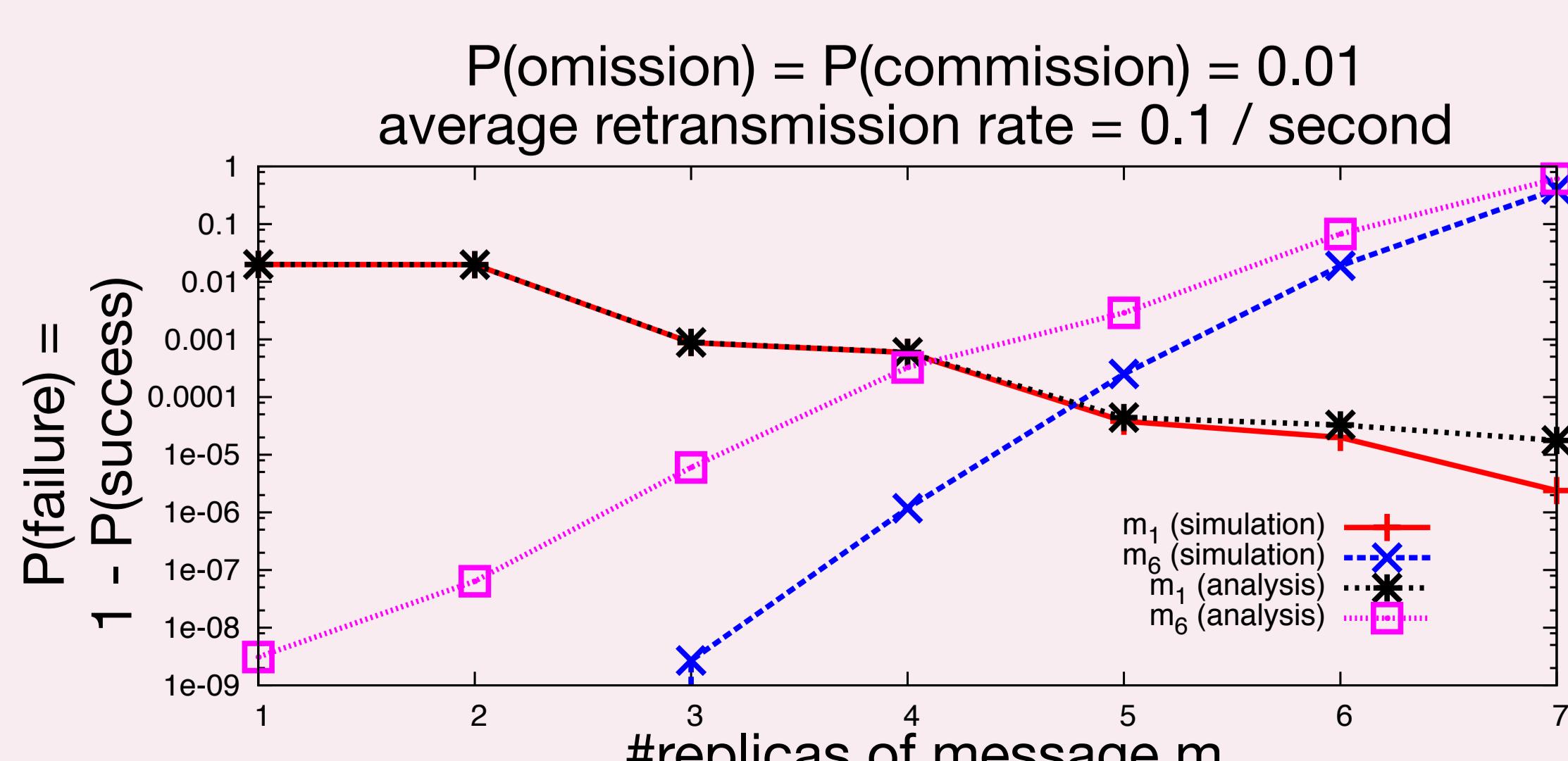


Probabilistic analysis

Fault model

- Poisson distribution for retransmissions
- Fixed host-specific probabilities for message omissions and message commissions

Simulation vs. analysis



Probability of “success” for single round of message transmissions

1. Compute probability that certain messages are omitted
2. For messages that were not omitted, compute probability of timeliness, inspired by Broster et al.’s method [1]
3. For messages that were not omitted and received on time, compute probability that a “majority” of messages are not corrupted
4. Repeat for all possible subsets of messages from this round

$$P_{su}(M_j) = \sum_{O_j \subseteq M_j}^4 P_{om}(O_j) \cdot \sum_{k \leq |O'_j|}^1 P_{ti}(O'_j, k) \cdot P_{co}(O'_j, k)$$

[1] I. Broster, A. Burns, and G. Rodríguez-Navas, “Timing analysis of real-time communication under electromagnetic interference,” Real-Time Systems, 2005.