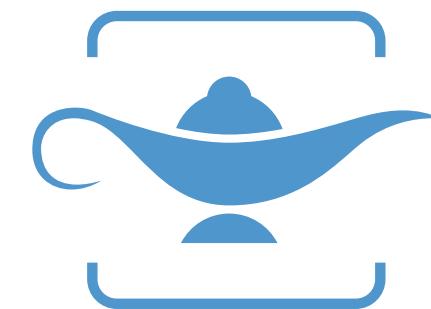


From Iteration to System Failure

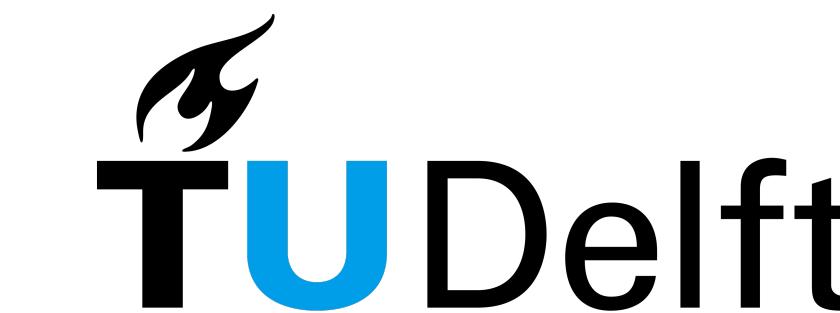
Characterizing the FITness of Periodic Weakly-Hard Systems

Arpan Gujarati*, Mitra Nasri#, Rupak Majumdar*, and Björn B. Brandenburg*

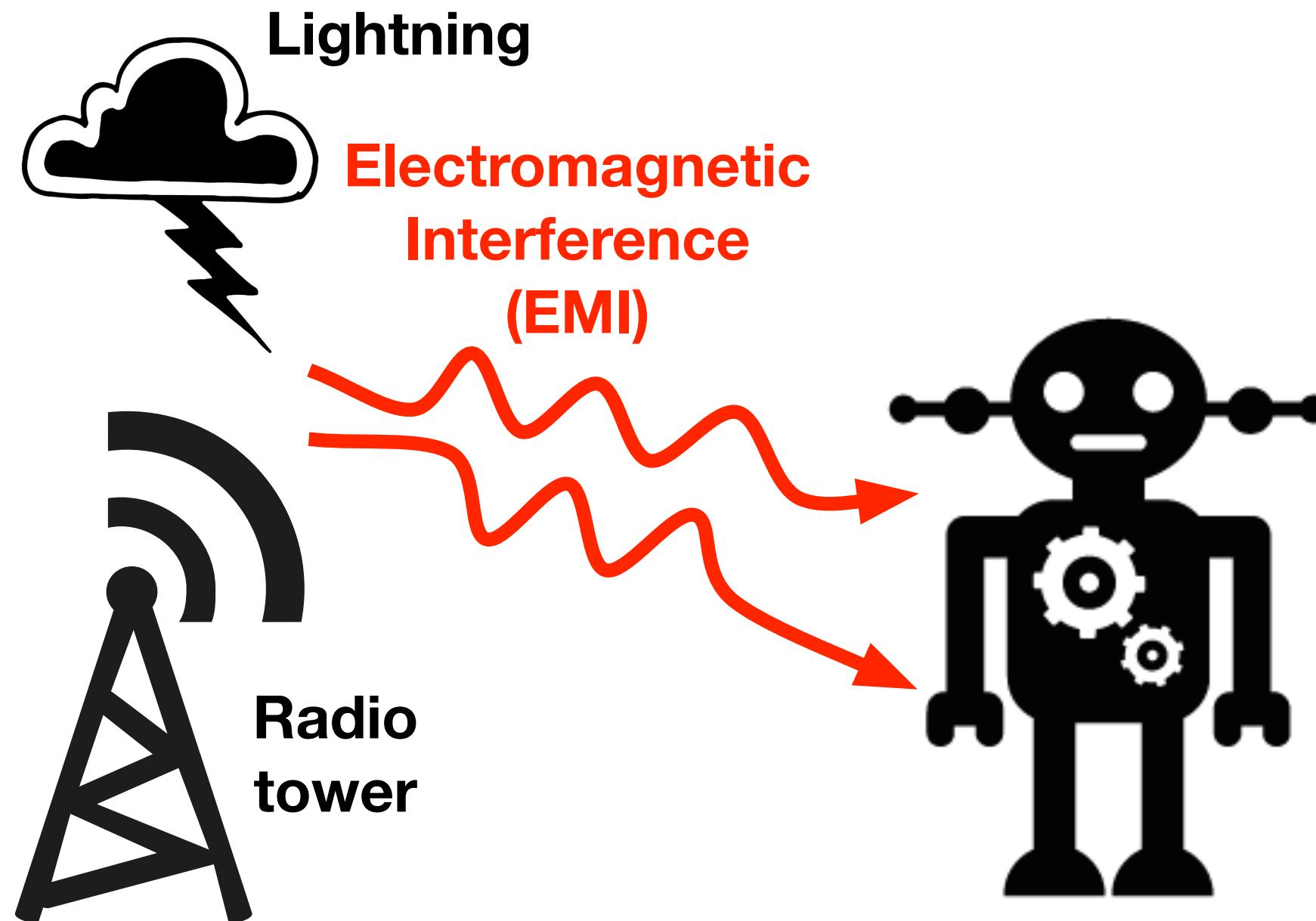
*MPI-SWS (Germany), #TU Delft (Netherlands)



MAX PLANCK INSTITUTE
FOR SOFTWARE SYSTEMS



Quantitative Reliability Analysis is Essential for Safety-Critical CPS



Zero risk of failures can never be achieved

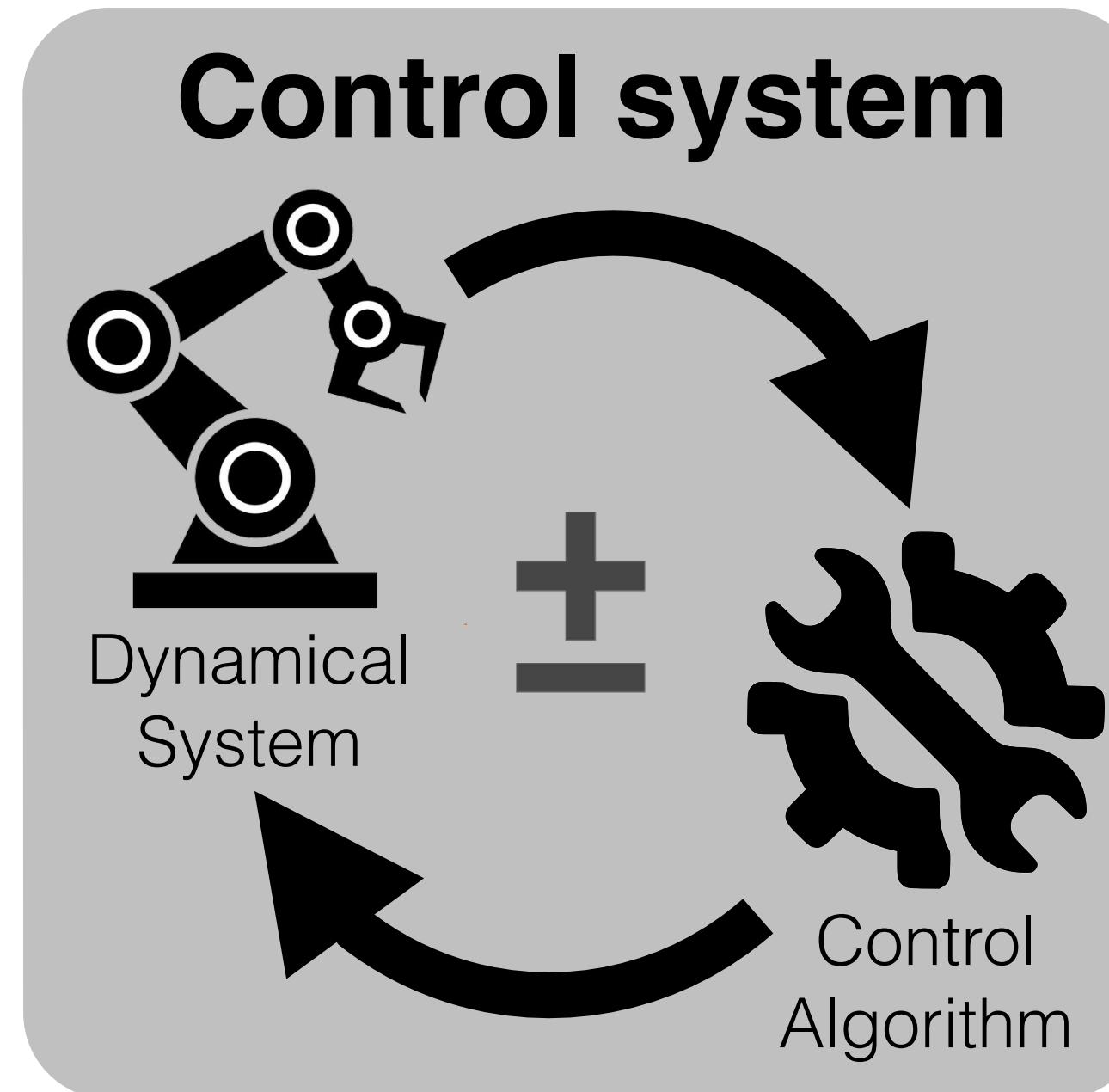
Safety certification objective:
Ensure “**negligible**” failure rates



E.g., for critical subsystems:
 $\text{Pr}[\text{failure / hour}] < 10^{-9}$

ARP4761 and the Safety Assessment Process for
Civil Airborne Systems

How to Analyse the Reliability of Temporally Robust Systems?



Motivating example

- **Frequency:** 100 Hz (10 ms time period)
- **Stability requirement:** 3 out of 4 iterations execute on time
- **Schedulability analyses:** $\Pr[\text{single iteration delayed}] \leq 10^{-10}$

Per-iteration analyses yield pessimistic failure rates

- Computing mean time to first failed iteration ignores stability requirements
- E.g., iteration failure probability of $10^{-10} \rightarrow 36,000 \times 10^{-9}$ failures / hours

9 orders of magnitude!

Explicitly accounting for the stability requirements

- Yields more accurate failure rates
- E.g., iteration failure probability of 10^{-10} and stability requirement $\rightarrow 1.08 \times 10^{-15}$ failures / hours

Not trivial anymore!

This work

How to Analyse the Reliability of **Temporally Robust Systems**?

Objectives

Generic

- Complex robustness requirements

Accurate (ideally, exact)

- Minimize pessimism in the final system reliability

Scalable

- Asymptotic requirements with large parameter values

Proposed Techniques

PMC (Probabilistic Model Checking)

- Exact, very generic, but slow

Mart (uses martingale theory)

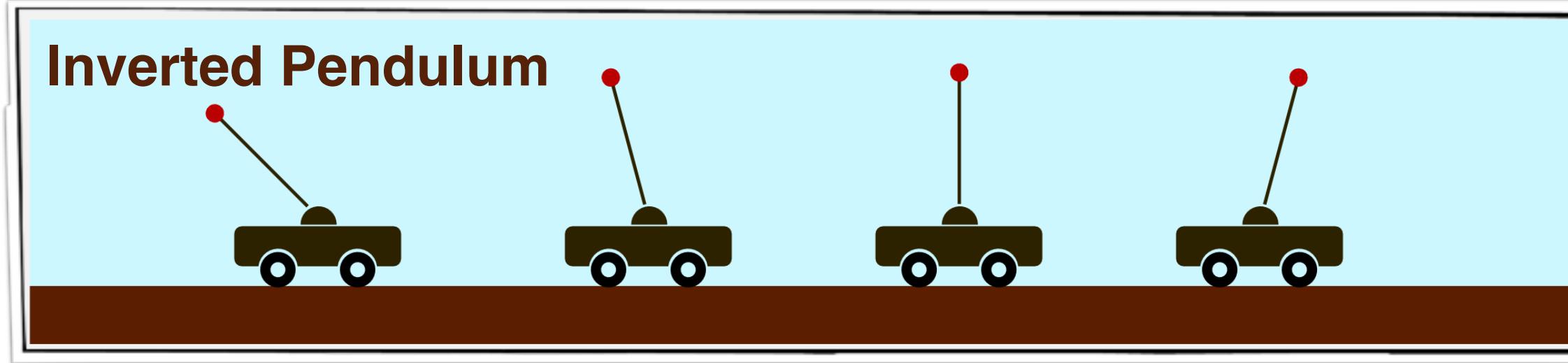
- Exact, less generic, but slightly faster

SAP (Sound Approximation)

- Not exact, least generic, but highly scalable

Background & System Model

Asymptotic Properties



Specification: Mass 0.5 kg, length 0.20 m, period 10 ms

Design: Current iteration is skipped \mapsto Use previous iteration parameters

Asymptotically stable with **at least 76.51% successful iterations***



Doesn't specify if the system can handle a burst of skipped iterations

→ What if the first 50 iterations are skipped? No feedback for 0.5 second

* Majumdar et al. "Performance-aware scheduler synthesis for control systems." EMSOFT, Taipei (2011)

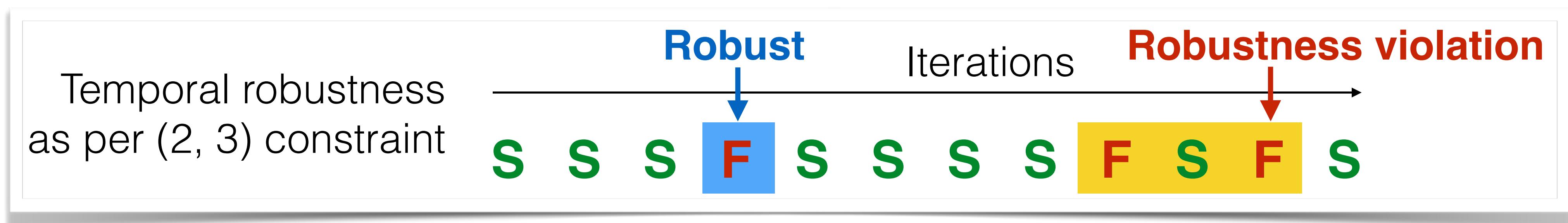
Weakly-Hard* Constraints



Concretize asymptotic properties using finite window sizes

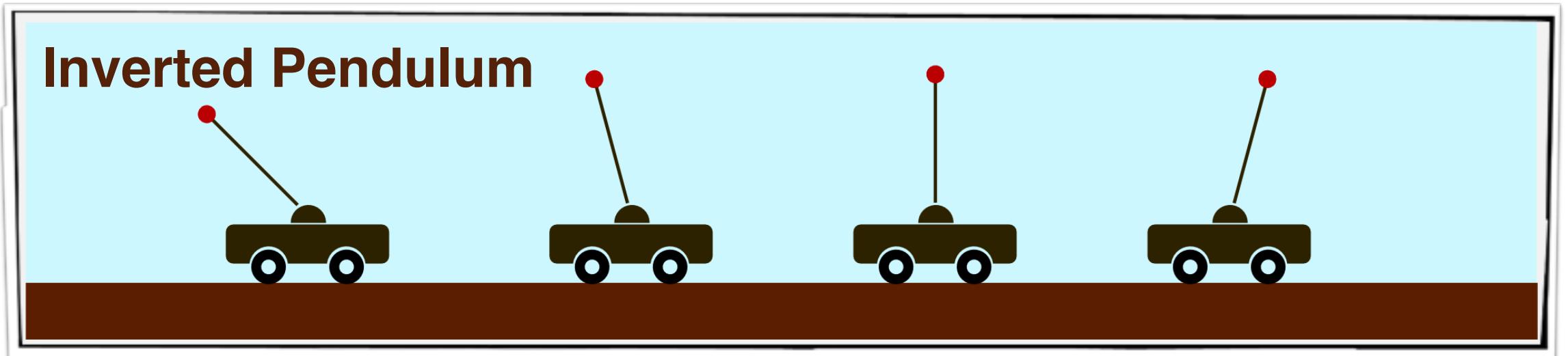
If each iteration is labeled either as a **Success** or a **Failure**

→ (m, k) constraint: At least m out of every k consecutive iterations must be **Successful**



* Bernat et al. "Weakly hard real-time systems." *IEEE Transactions on Computers*, 50(4):308–321 (2001).

Temporal Robustness Criteria



Asymptotically stable with **at least 76.51% successful iterations***

Short-range “liveness” constraints

- The inverted pendulum can tolerate a small burst of skipped iterations

Robustness Criteria

Combination of two weakly-hard constraints

(766, 1000)

(1, 5)

Combination of different weakly-hard constraints

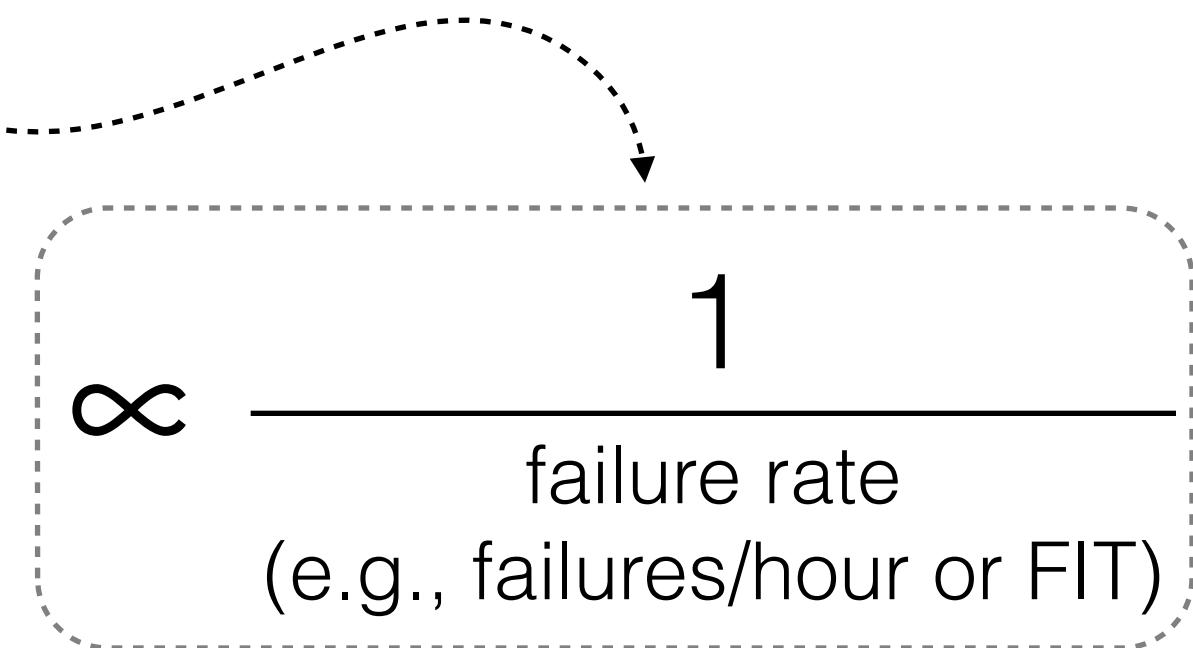
- (m, k) = Each k consecutive iterations, at least m successes needed
- $\langle m, k \rangle$ = Each k consecutive iterations, at least m consecutive successes needed
- $\overline{\langle m \rangle}$ = m consecutive failures should never happen

Problem Statement

Given periodic system **S**, time period **T**, iteration failure probability **P_F**, and the **temporal robustness criteria** ...

Lower-bound the **Mean Time To Failure (MTTF)** of S

$$\begin{aligned} \text{MTTF} &= \text{Expected time to 1st temporal robustness violation} \\ &= \sum_{n=0}^{\infty} \left(nT \times \Pr[1^{\text{st}} \text{ violation in the } n^{\text{th}} \text{ iteration}] \right) \end{aligned}$$



Assumption: **P_F** is independently and identically distributed (IID)^{1, 2}

¹ Broster et al. "Timing Analysis of Real-Time Communication Under Electromagnetic Interference." Real Time Systems Journal (2005)

² Gujarati et al. "Quantifying the Resiliency of Fail-Operational Real-Time Networked Control Systems." ECRTS, Barcelona (2018)

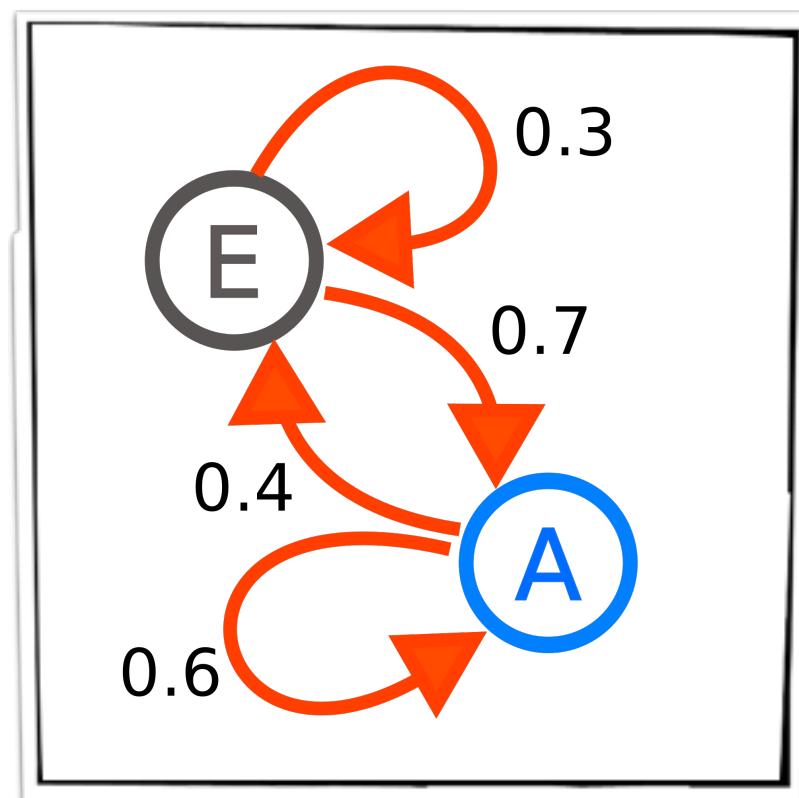
Probabilistic Model Checking (PMC)

Exact, very generic, but slow

MTTF Estimation using PMC

Periodic system **S** with iteration failure probabilities **P_F**

System as a probabilistic model



Safety properties as temporal logical specifications

$R=? \quad [\text{F safe}=\text{false}]$

Violation of temporal robustness

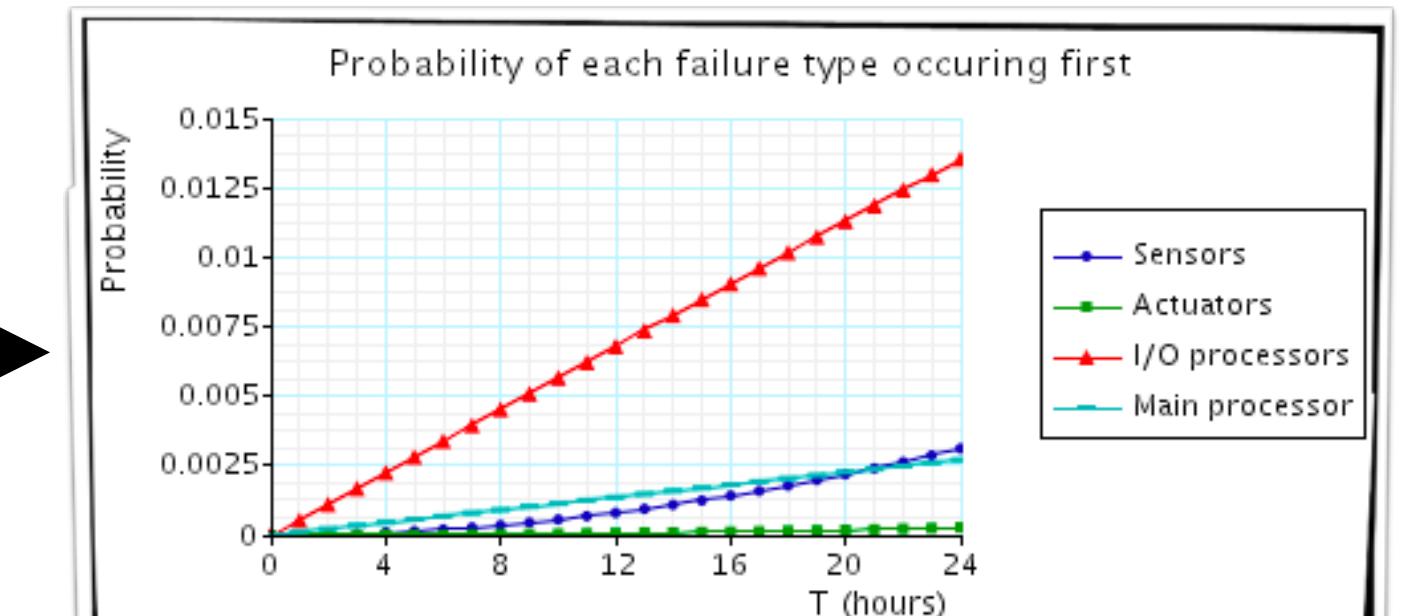
Formal verification technique to model and analyze systems that exhibit **probabilistic** behaviours

Probabilistic model checker, e.g., PRISM

```
PRISM 3.1
File Edit Model Properties Options
PRISM Model File: /home/.../power_policy1.sm
Model: power_policy1.sm
Type: Stochastic (CTMC)
Modules:
  SQ:
    q: min: 0; max: q_max; init: 0;
  SP:
    sp: min: 0; max: 2; init: 0;
  PM
Constants:
  q_max: int;
  rate_arrive: double;
  rate_serve: double;
  rate_s21: double;
  rate_i2s: double;
  q_trigger: int;
  ...
  // Service Queue (SQ)
  // Stores requests which arrive into the system to be processed.
  const int q_max = 20;
  const double rate_arrive = 1/0.72; // (mean inter-arrival time is 0.72 seconds)
  module SQ
    // q = number of requests currently in queue
    q: [0..q_max] init 0;
    // A request arrives
    [request] true -> rate_arrive : (q'=min(q+1,q_max));
    // A request is served
    [serve] q>1 -> (q=q-1);
    // Last request is served
    [serve_last] q=1 -> (q=q-1);
  endmodule
  ...
  // Service Provider (SP)
  // Processes requests from service queue.
  // The SP has 3 power states: sleep, idle and busy
  const double rate_serve = 1/0.008;
  const double rate_s21 = 1/1.6;
  const double rate_i2s = 1/0.67;
  ...
  // Rate of service (average service time = 0.008s)
  // Rate of switching from sleep to idle (average transition time = 1.6s)
  // Rate of switching from idle to sleep (average transition time = 0.67s)
```

MTTF estimation using PRISM

Quantitative results

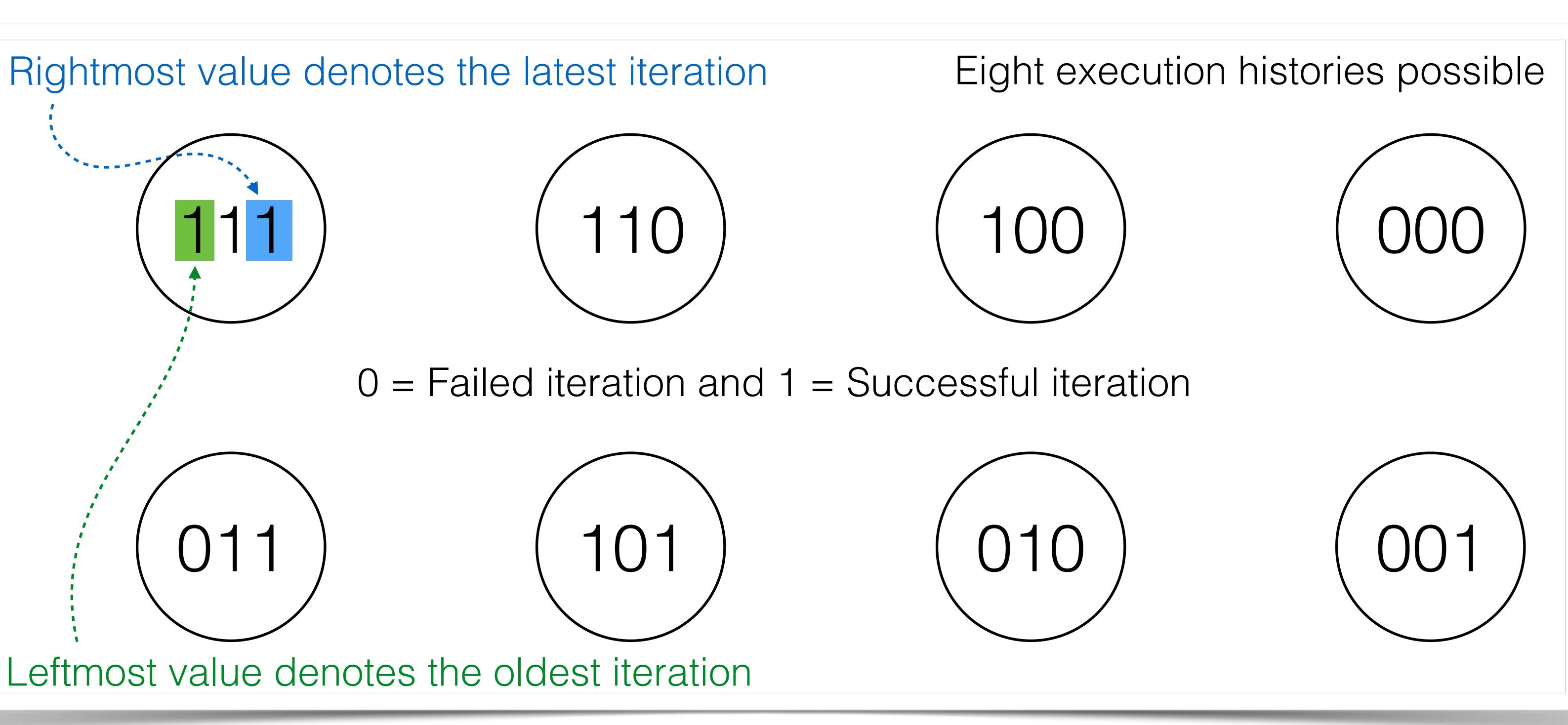


Modeling Weakly-Hard Constraints

Weakly-hard constraints depend on a **finite-sized history**

Key idea

- E.g., (m, k) constraint depends on the k latest iterations
- Connect all possible execution histories via transition probabilities P_F and $1 - P_F$



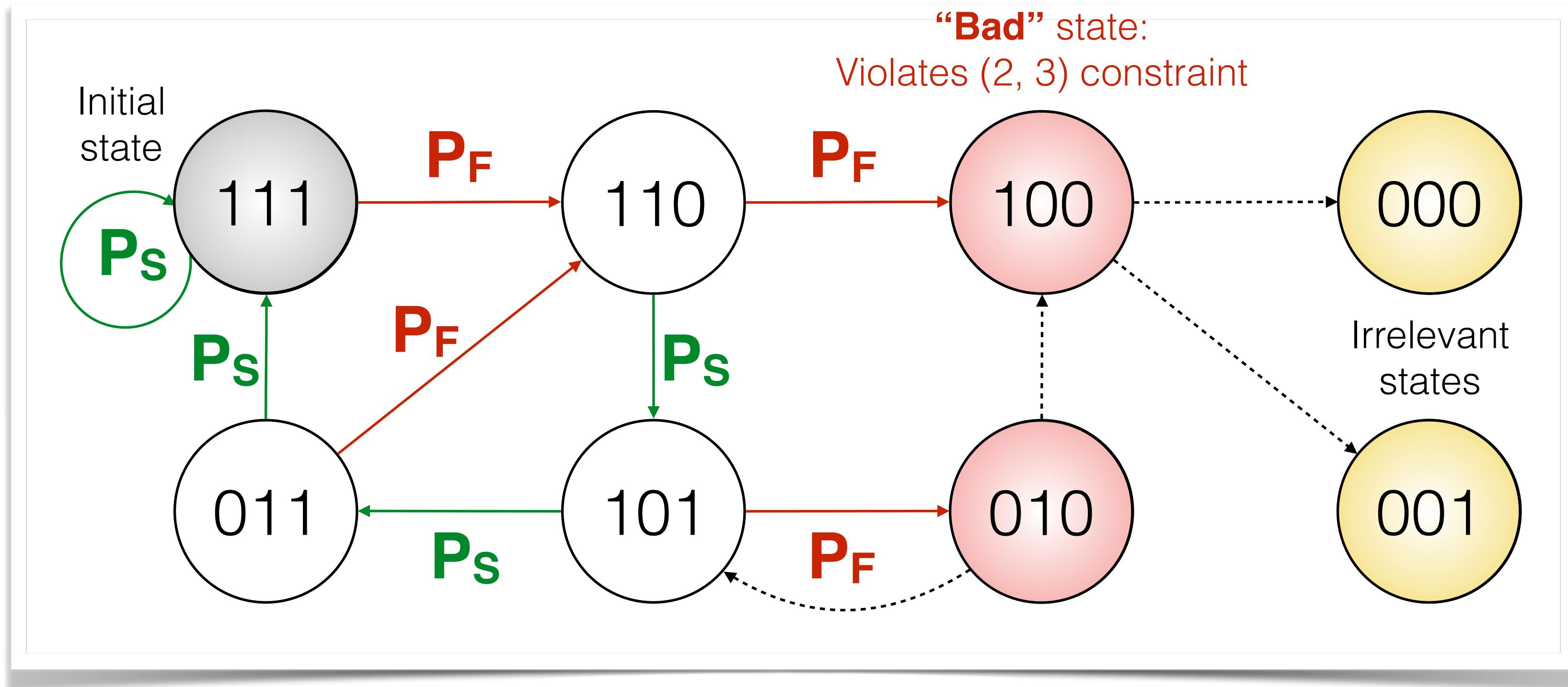
Example:
(2, 3) constraints

Modeling Weakly-Hard Constraints

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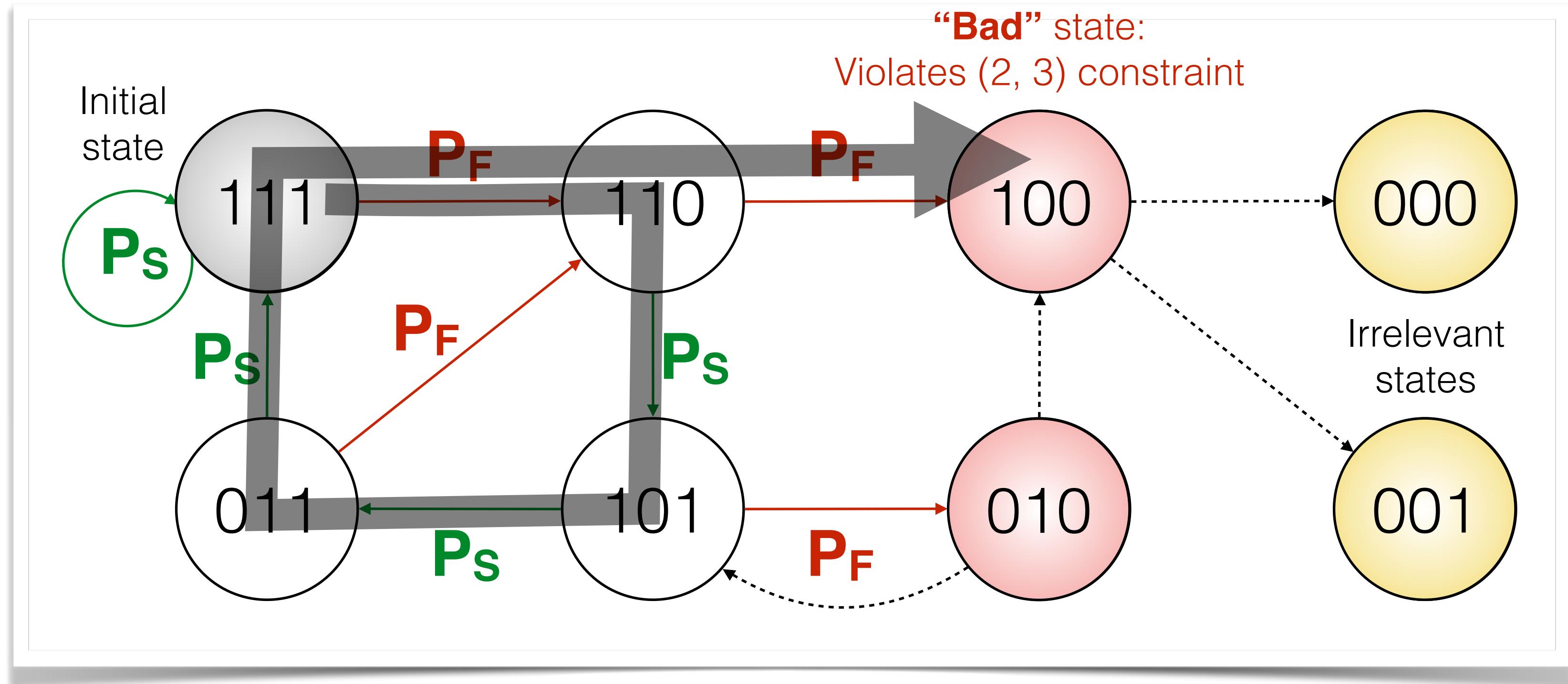
Example:
 $(2, 3)$ constraints

Modeling Weakly-Hard Constraints

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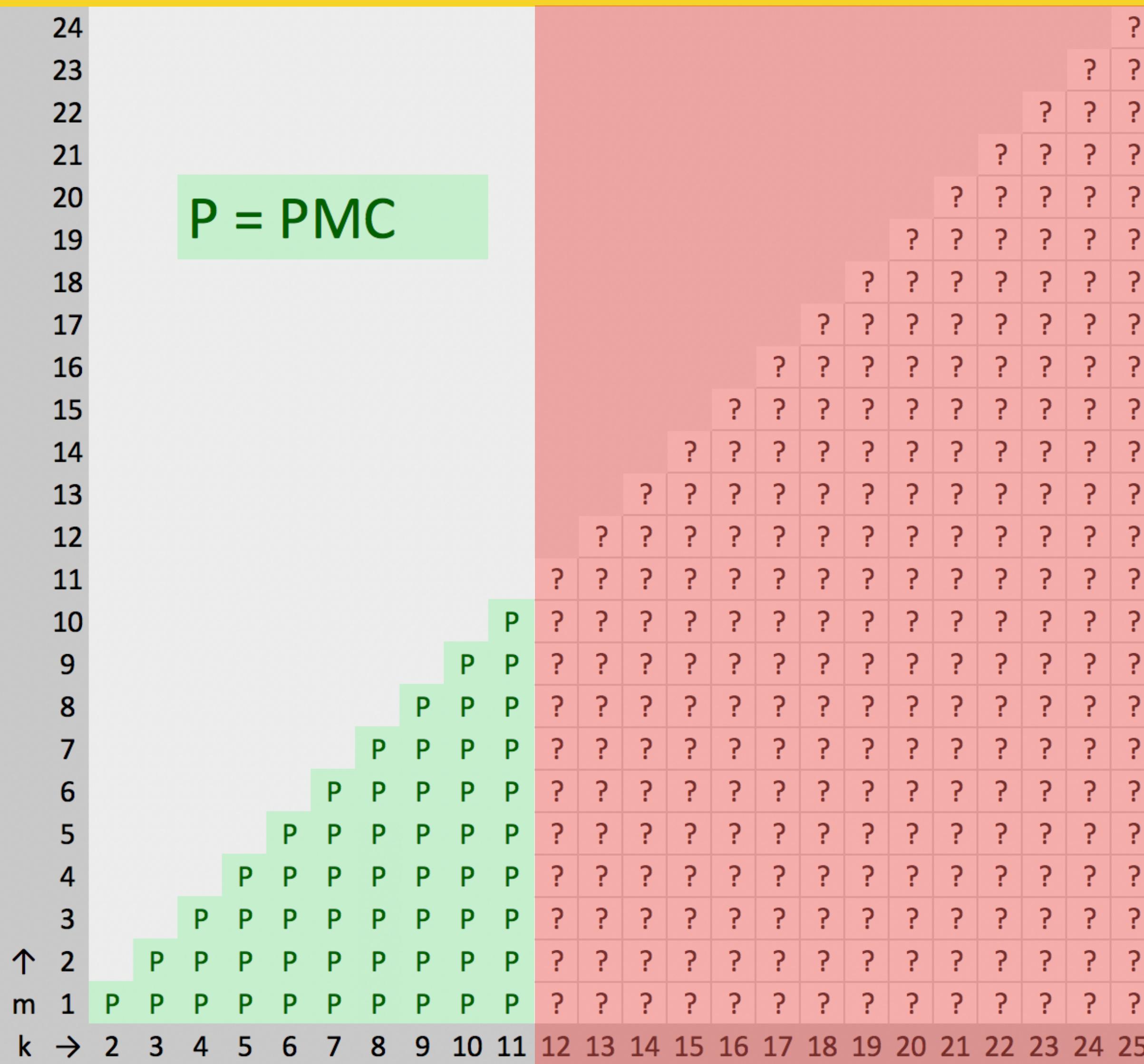


Example: (2, 3) constraints

$$\text{MTTF} = \left(\begin{array}{l} \text{Expected \# steps} \\ \text{to a “bad” state} \end{array} \right) \times T$$

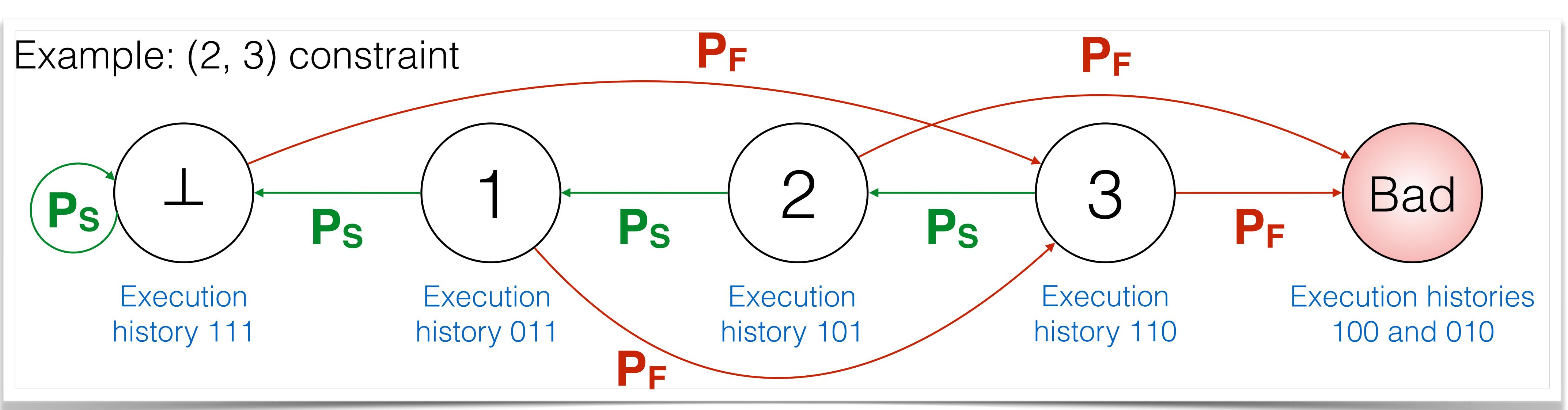
PRISM

Does PMC Scale with k?

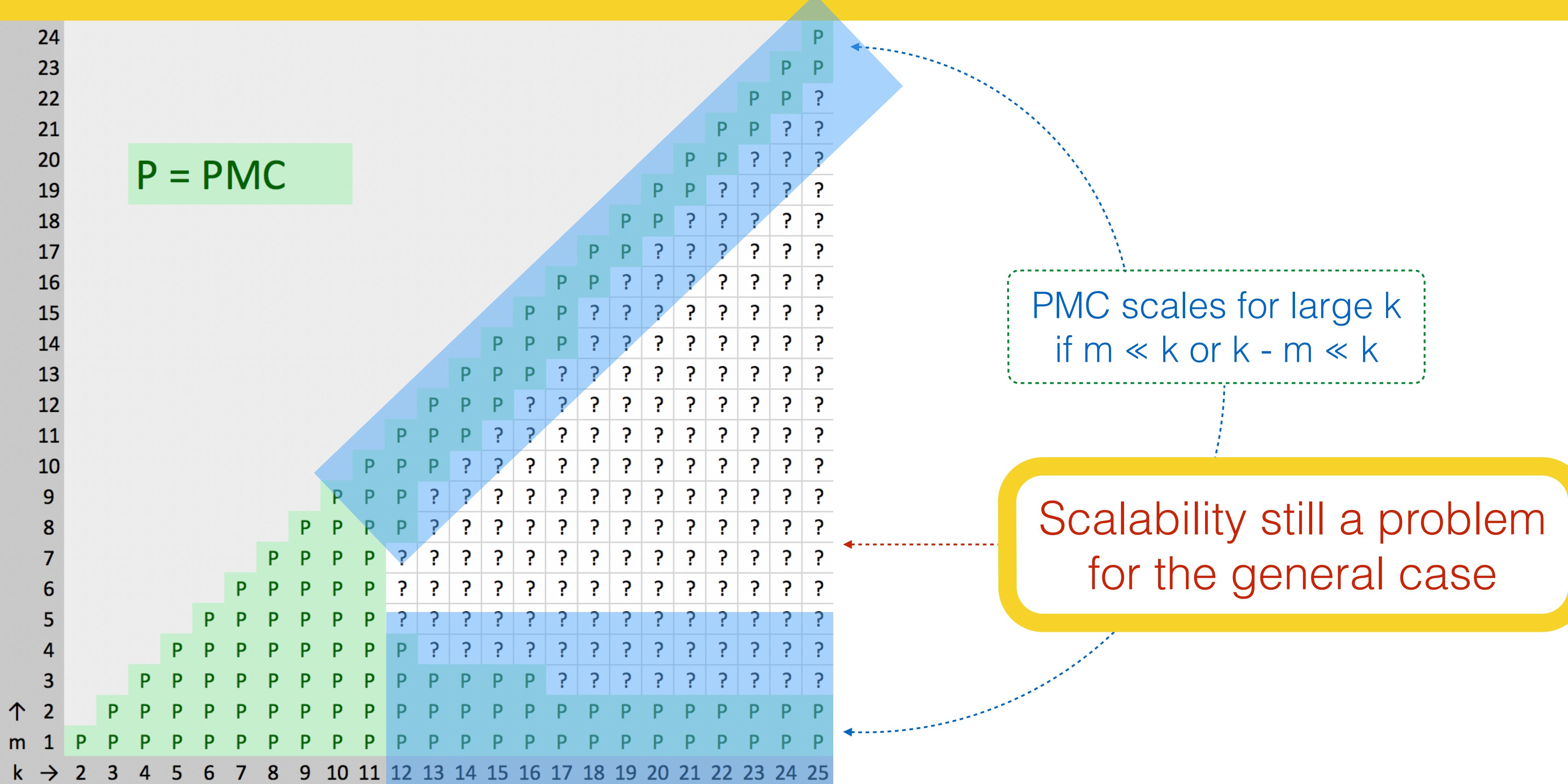


Optimizing for the Common Case $k - m \ll k$

Store **positions of all failed iterations**, instead of the entire history



Does the Optimized PMC Scale with k?

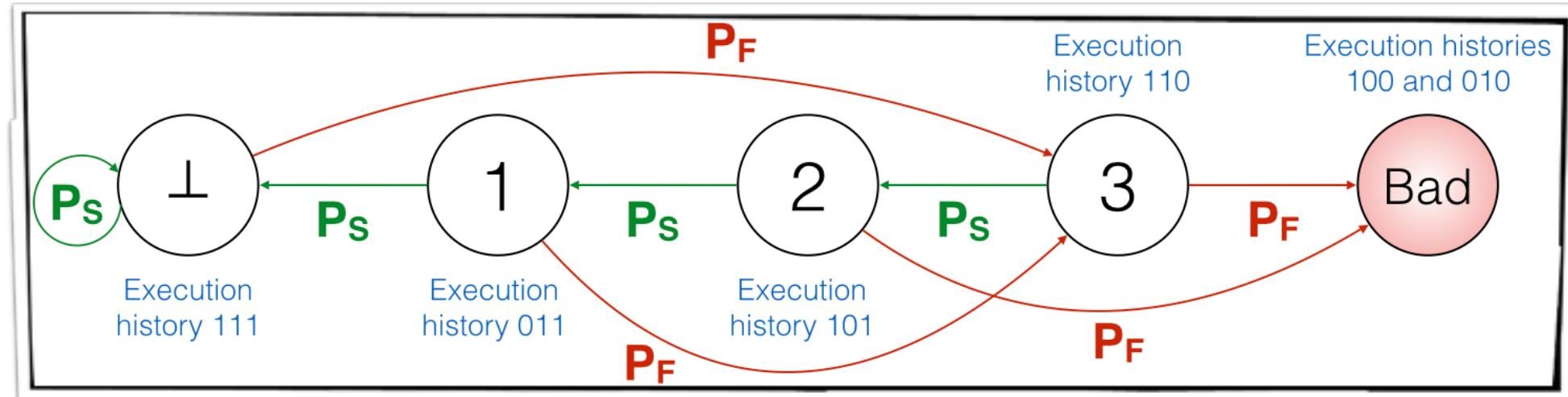


The Martingale Approach (**Mart**)

Exact, less generic, but slightly faster

Exact Model Checking Slows Down PRISM

Markov model



System of linear equations

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ -1 & 110 & 10 & 110 \\ -1 & 10 & 1090/9 & 10 \\ -1 & 0 & 100/9 & 1000/9 \end{bmatrix} \begin{bmatrix} e_0 \\ e_1 \\ e_2 \\ e_3 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

Model building

Model solving

MTTF (expected reward)

Probabilistic model checking
(PRISM under the hood)

For error-free
computation

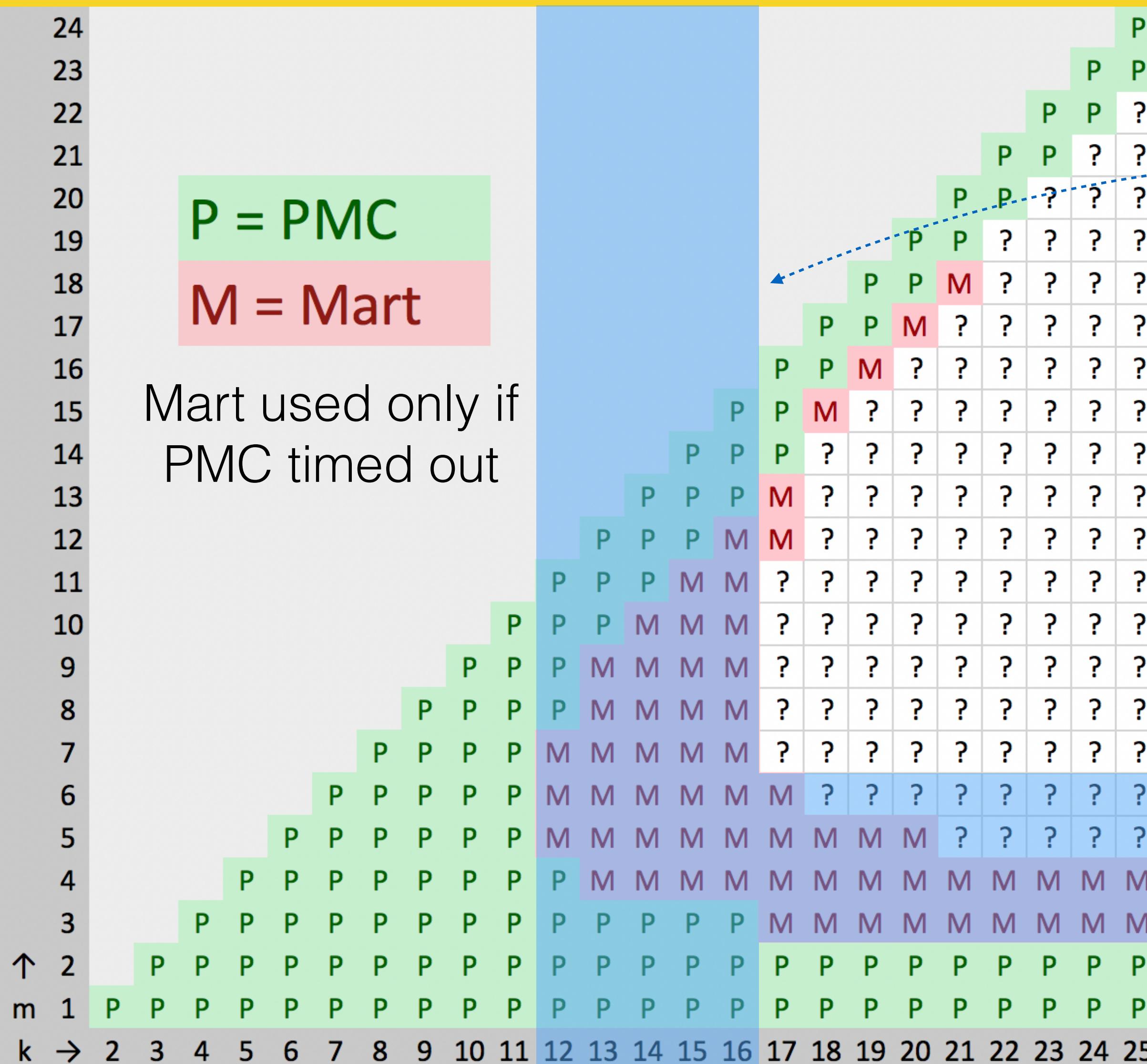
PRISM must be configured
with **exact model checking**
(i.e., no floating points)

Using **martingale theory***

- Linear equations obtained directly
- Bypass PRISM, use highly-scalable BLAS/LAPACK libraries, with very high precision

* Li. "A Martingale Approach to the Study of Occurrence of Sequence Patterns in Repeated Experiments." The Annals of Probability 8.6 (1980):1171–1176.

Mart Scales Better than PMC



Mart helps scale up exact MTTF estimation to **k = 16**

Also, Mart implicitly benefits from small values of m

Scalability still a problem for the general case

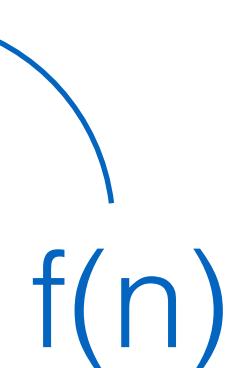
Sound Approximation (SAp)

Not exact, least generic, but highly scalable

Sound Approximation (SAP) for Single (m, k) Constraint

MTTF = Expected time to 1st temporal robustness violation

$$= \sum_{n=0}^{\infty} (nT \times \Pr[1^{\text{st}} \text{ violation in the } n^{\text{th}} \text{ iteration}])$$



Approximation accuracy

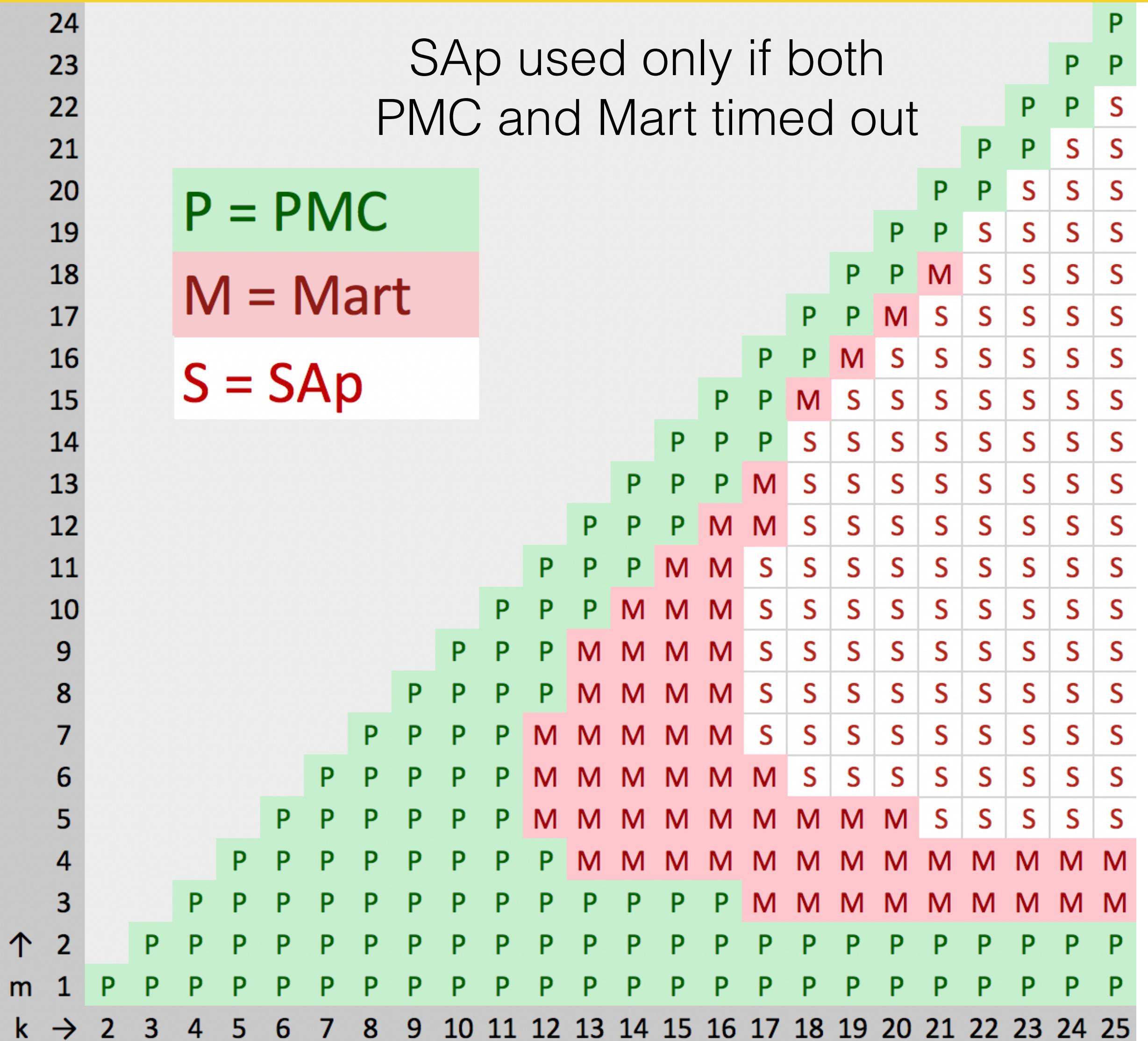
- Accuracy of $f_{LB}(n)$ (reliability modeling literature*)
- The choice of $n_0, n_1, n_2, \dots, n_D$ (heuristics based on $f_{LB}(n)$'s shape)

MTTF_{LB}

- ① Obtain $f_{LB}(n) \leq f(n)$ that can be quickly computed for large n
- ② Compute $f_{LB}(n_0), f_{LB}(n_1), \dots, f_{LB}(n_D)$
- ③ **Numerically integrate** over subintervals $(n_0, n_1], \dots, (n_{D-1}, n_D]$

* Sfakianakis et al.. "Reliability of a consecutive k-out-of-r-from-n: F system." IEEE Transactions on Reliability 41.3 (1992): 442-447.

SAp is Scalable to Very Large Window Sizes



SAp comfortably scales for
windows of size $k = 1000$

How Accurate is SAp?

All errors are positive (SAp is proven to under-approximate the exact MTTF)

11										81.45
10	Error > 50 %									
9	25 % < Error ≤ 50 %									79.25 69.20
8	Error ≤ 25 %									76.61 65.85 58.25
7										73.38 62.01 54.73 47.79
6										69.36 57.75 50.31 44.30 39.06
5										64.29 52.99 44.91 40.11 35.35 31.43
4										57.84 47.34 39.12 34.28 31.22 28.13 25.06
3										49.62 39.96 33.44 28.17 24.82 22.98 21.64 20.28
2										39.29 30.21 25.73 22.68 20.09 17.80 15.93 14.55 13.63
m	↑ 2	25.91	19.44	15.77	13.60	12.30	11.50	10.98	10.62	10.35 10.13
1	05.76	05.76	05.76	05.76	05.76	05.76	05.76	05.76	05.76	05.76
k →	2	3	4	5	6	7	8	9	10	11 12

SAp is reasonably accurate

Example: If $\text{MTTF}_{\text{exact}} = 10^9$ hours,
100% error $\rightarrow \text{MTTF}_{\text{SAp}} = 0.5 \times 10^9$ hours

Relative errors significant even for small k

- Exact analysis needed when feasible

Summary

Approach	Accuracy	Expressiveness	Scalability
PMC	Exact	General system, any weakly-hard constraint	Poor ($k \leq 11$)
Mart	Exact	IID systems, any weakly-hard constraint	Poor ($k \leq 16$)
SAp	Sound approx. ($\leq 100\%$)	IID systems, single (m, k) constraint	Good ($k \leq 1000$)

Future work: Make **SAp more expressive**

- Handle other / multiple weakly-hard constraints
- Beyond IID iteration failure probabilities

More in the paper!

- PRISM code and Mart example
- PMC / Mart for $\langle m, k \rangle$ and $\langle \overline{m} \rangle$ constraints
- SAp details and soundness proofs
- More extensive evaluation of PRISM