

EIG Algorithm for Byzantine Agreement



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Exponential Information Gathering for Byzantine Agreement

Reaching consensus is extremely important in \$1000 \$2000 any distributed network.

eg: we cannot have two datanodes in a cluster such that one thinks price = \$1000 while the other thinks price = \$2000

Depending on which node the stequest hits. The user would see the corresponding value, giving an inconsistent view

Somehow, the nodes need to agree on one value.

Byzantine Agreement

Byzantine agreement is a problem of

reaching a consensus even when one or

many nodes processes are malicious corrupt



Core Idea: Relay the values across rounds, record the communication path, and decide.

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EIG Data Structure

Fix data structure is a tree that grows exponentially. The paths from the root of the tree represent the communication path from which the message is received (propagaked)

The tree is constructed level by level and is designed to hold all possible permutations of length k. (distinct paths)

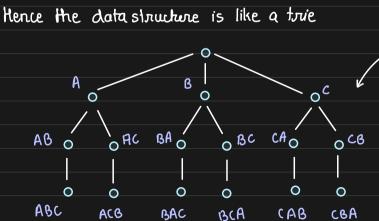
At each level k, every node has

Root node is labelled as ""

if a node succeived a message labelled [3,2,7,1,4]

n-k children to maintain uniqueness of the path

it holds the message in the tree along that path



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3-level deep EIG

constructed over 3

nodes A,B and C

The algorithm

* construction of Elle Tree is covered in previous topics

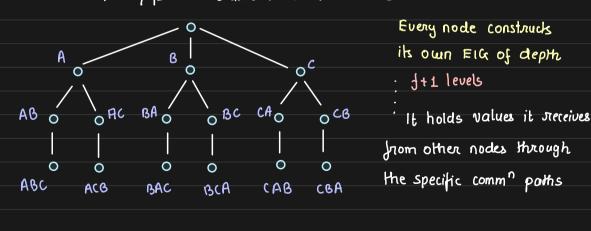
we assume EIG Tree is independently constructed at each node

The algorithm is tolerant to f faulty processes solong as n > 3f i.e number of nodes are much larger than the faults

When a process sends "ill-formed" (junk/garbage) value to other nodes participating in the consensus, the consensus should not be prone to that.

The How

The processes propagaks values for f+1 rounds and each node process builds its own EIG Tree



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if a node steceives garbage value, the node ignores it eg: value should be into but string received value should be in range [0-100], but 5000 received Once the Elk is constructed to make the decision, a node traverses the tree bottom-up The value from the leaves, is propagated up for non-leaf node, value = majority of values 1 of its children if majority, vz = maj(v) 0 else, vn = vo ACB BAC BCA CAB CBA the default value

The final consensus value will be value converged at root.

1. if no faulty nodes, no corrupt nodes, all nodes would converge to the same $v_{\varkappa} = v$

2 few counupt nodes, sending counupt values will be absorbed and not propagated up.

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