

# Synch BFS in Distributed Systems



## Synchronous BFS in distributed systems

Breadth-First Search is a critical algorithm in Distributed Systems because it powers some key features like

- Broadcast in Minimum Hme
- Building topologial understanding
- Topological Stat like Diameter and Shortest path across nodes

We discuss a Synchronous BFS algorithm in which all nodes synchronize and move ahead with each round in sync.

Output of BFS

The output of this traversal is a

breadh-first directed spanning tree

all nodes covered with min edges

Each node knows its parent and its children

\* Node may be reached via multiple nodes but in BFS spanning tree, only one will be the panent

The algorithm Because the algorithm is synchronous, nodes proceed in sync Say, node in initiates the BFS and there are total n nodes. Round 1: is sends search message to its neighbours When an unmarked node receives the message - it manks itself to node it received message - updates its parent from Round 2: the nodes who received message in round 1 sends search message to its out going neighbours The steceiving nodes takes the action by marking themselves and continuing the process.

Eventually every single node will be covered

and we would be having a direcked spanning tree

Complexity Analysis

Given that the nodes proceed every round, the time taken to

cover the entire network will be proportional to the farthest

node & diameter of the network

The # messages exchanged ie. communication complexity will be at max the number of edges in the network.

With our BFS implementation, every node upon steeliving an incoming search message, knows its parent, but how would a parent know its children?

\* Outgoing edge & child we are building a spanning tree here

An node, upon stectiving the search message, would if unmarked, updale the local state with parent

and inform back if the node is chosen as

if marked, discard the message

parent or non-parent

Conveying child pointers

How would our algorithm terminate?

The most impositant pasit of any distributed algarithm is its termination.

How would the node know that BFS is done?

The approach we use is called Converge cast

Cone idea: what if the nodes mespond to their panent only after they got mesponse from their descending nodes

Kind of like Post order traversal but in a distributed setup.

Nodes would respond back with parent non-parent messages Instead of responding immediately, and we wait until we get response from descendents

Thus, the most node will get nesponse

only after the entire network is covered. This is called Convergecast

If edges are bi-directional, then above approach warks

if they are uni-directional, then each node would brigger another BFS to inform the parent

Applications Broadcast: If a node in a distributed system wants to broad cast a message m, it can initiale a Synch BFS cuith itself as the root. The message will be propagated throughout the tree from parent to children. Distributed computation: In order to use the computation capability of the entire network the computation is sent from root to all the nodes and the results are propagated back

Computing diameter of the network:

Each node initiates a BFS and thus knowing the farthest node. The max distances observed by each node is then flooded through another BFS to compute the global max ie. diameter.

to the root - like a post order traversal