







Complete and Efficient DRAT Proof Checking

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DRAT proof sequence of clause introductions and deletions

- i: C C must be a RUP or RAT clause
- d: C always allowed
- i: ☐ finishes the proof



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DRAT checker checks that each inference is correct



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RUP / RAT introduction and clause deletion are satisfiabilty-preserving



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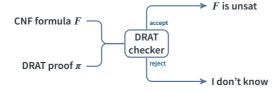
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DRAT checker checks that each inference is correct

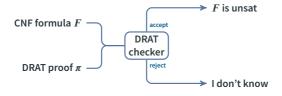
RUP / RAT introduction and clause deletion are satisfiabilty-preserving if \Box is derived, then the input formula is unsatisfiable

DRAT checker checks that each inference is correct

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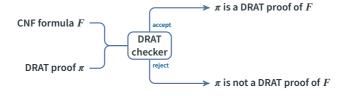


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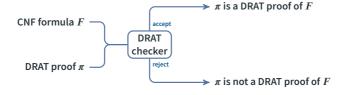
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Problem state-of-the-art DRAT checkers decide something slightly different from DRAT correctness

Proof system specification of *correct* refutations

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Sound proof system if a correct refutation of F exists, then F is unsatisfiable

Refutations are correct with respect to a proof system

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specified DRAT



[Wetzler, Heule, Hunt '14]

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drat-trim, gratgen

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drat-trim, gratgen
unit deletions are ignored
sound proof system

Discrepancies exist [Rebola-Pardo, Biere '18] sDRAT-correct proofs that are oDRAT-incorrect, and oDRAT-correct proofs that are sDRAT-incorrect

Why operational DRAT? The methods used in state-of-the-art DRAT checkers do not work when unit clause deletions occur in the input proof ignore unit clause deletions

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Contributions

- Can we check sDRAT as efficiently as oDRAT?
- How often do discrepancies occur in practice?

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- How often do discrepancies occur in practice? Quite a lot

Experimental evaluation

Questions

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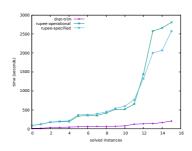
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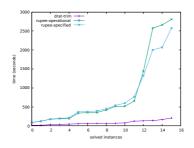
Benchmarks 11 instances from SAT Competition 2017

Solvers COMinisatPS_Pulsar,glucose-4.1,Maple_LCM_Dist,cadical-sc17

Experimental results

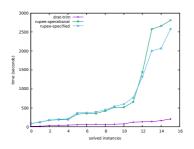


Experimental results



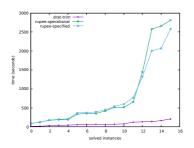
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Experimental results



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rupee in sDRAT mode performs comparably to rupee in oDRAT mode no observable overhead due to unit clause deletion

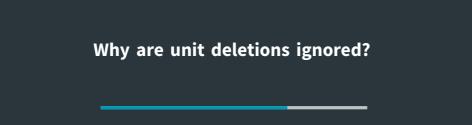


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All proofs generated by cadical-sc17 are correct for both sDRAT and oDRAT. For the other 3 solvers, 8 out of 11 proofs were incorrect as sDRAT and correct as oDRAT.

probably due to the CDCL proof generation method inherited from minisat



State-of-the-art DRAT checking

x_1	$x_5 x_6$	$\overline{x_3} \overline{x_6} x_8$	$\overline{x_4}\overline{x_9}x_{10}$
$\overline{x_1} x_2$	$\overline{x_2}\overline{x_5}x_7$	$x_3 \overline{x_4} \overline{x_6}$	$x_9 \overline{x_{10}}$
$\overline{x_1} \overline{x_2} x_3$	$\overline{x_1} \overline{x_5} x_6$	$x_5 \overline{x_8}$	$x_7 \overline{x_9}$
$\overline{x_1} \overline{x_3} x_4$	$x_4 \overline{x_5} \overline{x_6}$	$\overline{x_3} x_9 x_{10}$	$\overline{x_7} \overline{x_8} \overline{x_9} \overline{x_{10}}$

x_1	$x_5 x_6$	$\overline{x_3}\overline{x_6}x_8$	$\overline{x_4} \overline{x_9} x_{10}$
$\overline{x_1} x_2$	$\overline{x_2} \overline{x_5} x_7$	$x_3 \overline{x_4} \overline{x_6}$	$x_9 \overline{x_{10}}$
$\overline{x_1} \overline{x_2} x_3$	$\overline{x_1} \overline{x_5} x_6$	$x_5 \overline{x_8}$	$x_7 \overline{x_9}$
$\overline{x_1} \overline{x_3} x_4$	$x_4 \overline{x_5} \overline{x_6}$	$\overline{x_3} x_9 x_{10}$	$\overline{x_7} \overline{x_8} \overline{x_9} \overline{x_{10}}$

i: *x*₅

 $i: x_9$

i: □

i: *x*₅

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C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

 $(i: x_5)$

 $i: x_9$

i: □

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$$(i: x_5) \quad \overline{x_5}$$

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$$(\mathbf{i}: x_5)$$
 $\overline{x_5}$ x_1

i:
$$x_9$$

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$$(i: x_5) \quad \overline{x_5} \ x_1 \ x_2$$

$$i: x_0$$

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$$(\mathbf{i}: x_5) \quad \overline{x_5} \quad x_1 \quad x_2 \quad x_3$$

i:
$$x_9$$

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$$(i: x_5) \quad \overline{x_5} \quad x_1 \quad x_2 \quad x_3 \quad x_4$$

i: x_9

i: □

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

(i:
$$x_5$$
) $\overline{x_5}$ x_1 x_2 x_3 x_4 x_6

i: x_9

i: □

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$$(i: x_5) \quad \overline{x_5} \ x_1 \ x_2 \ x_3 \ x_4 \ x_6 \ x_8$$

i: x_9

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(i:
$$x_5$$
) $\overline{x_5}$ x_1 x_2 x_3 x_4 x_6 x_8 \perp

i: x_9

i: □

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$$(i: x_5) \quad \overline{x_5} \ x_1 \ x_2 \ x_3 \ x_4 \ x_6 \ x_8 \ \bot \ \Rightarrow RUP$$

i: x_9

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$$(\mathbf{i}: x_5) \quad \overline{x_5} \ x_1 \ x_2 \ x_3 \ x_4 \ x_6 \ x_8 \ \bot \ \Rightarrow \mathsf{RUP}$$

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$$\overline{\mathbf{i}}: x_9$$
 $\overline{x_9}$

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$$\overline{x_9} \quad \overline{x_9} \quad x_1$$

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i:
$$x_9$$
 $\overline{x_9}$ x_1 x_2

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$$(\mathbf{i}: x_9) \quad \overline{x_9} \quad x_1 \quad x_2 \quad x_3 \quad x_4$$

i: □

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$$(i: x_5) \quad \overline{x_5} \ x_1 \ x_2 \ x_3 \ x_4 \ x_6 \ x_8 \ \bot \ \Rightarrow RUP$$

i:
$$x_9$$
 $\overline{x_9}$ x_1 x_2 x_3 x_4 x_5

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 $\overline{x_9}$ x_1 x_2 x_3 x_4 x_5 x_6

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$$(i: x_5) \quad \overline{x_5} \ x_1 \ x_2 \ x_3 \ x_4 \ x_6 \ x_8 \ \bot \ \Rightarrow RUP$$

$$(i: x_9)$$
 $\overline{x_9}$ x_1 x_2 x_3 x_4 x_5 x_6 x_7

i: □

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(i:
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) $\overline{x_9}$ x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{10}

i: □

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$$(i: x_9) \quad \overline{x_9} \ x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6 \ x_7 \ x_8 \ x_{10} \ \bot$$

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$$(i: x_9)$$
 $\overline{x_9}$ x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{10} \bot \Rightarrow RUP

(i:
$$\square$$
) $x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6 \ x_7 \ x_8 \ x_9 \ x_{10} \ \bot \Rightarrow RUP$

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$$(i: x_5) \quad x_5 x_1 x_2 x_3 x_4 x_6 x_8 \perp \Rightarrow RUP$$

$$(i: x_9) \quad \overline{x_9} \ x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6 \ x_7 \ x_8 \ x_{10} \ \bot \ \Rightarrow \mathsf{RUP}$$

(i:
$$\Box$$
) $x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6 \ x_7 \ x_8 \ x_9 \ x_{10} \ \bot \Rightarrow RUP$

C is a reverse unit propagation (RUP) clause in F whenever $F \land \overline{C}$ reaches a conflict by unit propagation

$$\overbrace{\mathbf{i}: x_5} \quad \widehat{x_5} \quad x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_6 \quad x_8 \quad \bot \quad \Rightarrow \mathsf{RUP}$$

(i:
$$x_9$$
) $x_9 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{10} \bot \Rightarrow RUP$

$$(i: \Box)$$
 $x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_9 x_{10} \bot \Rightarrow RUP$

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$$(i: x_5) \quad \overrightarrow{x_5} \quad x_1 \quad x_2 \quad x_3 \quad x_4 \quad \overrightarrow{x_6} \quad x_8 \quad \bot \quad \Rightarrow \mathsf{RUP}$$

$$(i: x_9) \quad \overrightarrow{x_9} \quad x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_5 \quad x_6 \quad x_7 \quad x_8 \quad \overrightarrow{x_{10}} \quad \bot \quad \Rightarrow \mathsf{RUP}$$

$$(i: \Box)$$
 $x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_9 x_{10} \bot \Rightarrow RUP$

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

(i: □

i:
$$x_5$$
 x_5 x_1 x_2 x_3 x_4 x_6 x_8 $\bot \Rightarrow RUP$
i: x_9 x_9 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{10} $\bot \Rightarrow RUP$

 $x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6 \ x_7 \ x_8 \ x_9 \ x_{10} \perp \Rightarrow RUP$

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i:
$$x_5$$
 x_5 x_1 x_2 x_3 x_4 x_6 x_8 $\bot \Rightarrow RUP$
i: x_9 x_9 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{10} $\bot \Rightarrow RUP$

(i:
$$\square$$
) $x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_9 x_{10} \bot \Rightarrow RUP$

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

i:
$$x_5$$
 x_5 x_1 x_2 x_3 x_4 x_6 x_8 \bot \Rightarrow RUP
i: x_9 x_9 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{10} \bot \Rightarrow RUP
i: \Box x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{9} x_{10} \bot \Rightarrow RUP

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

i:
$$x_5$$
 x_5 x_1 x_2 x_3 x_4 x_6 x_8 $\bot \Rightarrow RUP$
i: x_9 x_9 x_1 x_2 x_3 x_4 x_5 x_6 x_7 x_8 x_{10} $\bot \Rightarrow RUP$

$$i: \square$$
 $x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6 \ x_7 \ x_8 \ x_9 \ x_{10} \ \bot \Rightarrow RUP$

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

 $[i: x_9]$

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

i: □

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

 $x_1 \ x_2 \ x_3 \ x_4 \ x_5 \ x_6 \ x_7 \ x_8 \ x_9 \ x_{10} \ \bot \Rightarrow RUP$

i: *x*₅

 $i: x_0$

i: □

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

Optimizations trace preprocessing and backwards checking

 x_1

i: *x*₅

i: x_0

i: □

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Optimizations trace preprocessing and backwards checking

 $x_1 x_2$

i: *x*₅

i: x_0

i: □

C is a reverse unit propagation (RUP) clause in F whenever $F \wedge \overline{C}$ reaches a conflict by unit propagation

Optimizations trace preprocessing and backwards checking

$$x_1 \ x_2 \ x_3$$

i: *x*₅

i: x_9

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7

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 x_0

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 $d: \overline{x_1} x_2$

i: *x*₉

i: □

 x_1 x_2 x_3 x_4

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Unit clause every literal but one is falsified by UP by the current formula reason clauses are always unit clauses

How do we delete unit clauses

i: *x*₅

$$x_1$$
 x_2 x_3 x_4 x_5 x_6 x_7 x_8

 $d: \overline{x_1} x_2$

i: x_9

i: □

i: □

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Unit clause deletion breaks stuff

$$x_{1} \qquad x_{5} x_{6} \qquad \overline{x_{3}} \overline{x_{6}} x_{8} \qquad \overline{x_{4}} \overline{x_{9}} x_{10} \qquad x_{5}$$

$$\overline{x_{1}} \overline{x_{2}} x_{3} \qquad \overline{x_{1}} \overline{x_{5}} x_{6} \qquad x_{5} \overline{x_{8}} \qquad x_{7} \overline{x_{9}}$$

$$\overline{x_{1}} \overline{x_{3}} x_{4} \qquad x_{4} \overline{x_{5}} \overline{x_{6}} \qquad \overline{x_{3}} x_{9} x_{10} \qquad \overline{x_{7}} \overline{x_{8}} \overline{x_{9}} \overline{x_{10}}$$

$$x_{1} \qquad x_{2} \qquad x_{3} \qquad x_{4}$$

$$i: x_{5}$$

$$x_{1} \qquad x_{2} \qquad x_{3} \qquad x_{4} \qquad x_{5} \qquad x_{6} \qquad x_{7} \qquad x_{8}$$

$$d: \overline{x_{1}} x_{2}$$

$$x_{1} \qquad x_{5} \qquad x_{6} \qquad x_{4} \qquad x_{3} \qquad x_{8}$$

$$i: x_{9} \qquad \text{How do we reconstruct the trace during backwards checking?}$$

$$i: \Box$$

Unit clause deletion breaks stuff

I'd wish.

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Two-watched literal schema makes unit propagation efficient

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Result negligible overhead of unit clause deletion!



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There seems to be no proof checking-based reason to prefer oDRAT over sDRAT.

Proof generation methods are inserting unnecessary deletions.

We conjecture that proofs can be generated according to the specification, and that proofs can be checked under the specification with similar performance to that of state-of-the-art DRAT checkers.