



String Sorting in Python - Comparison of Several Algorithms

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Comparison-based sorting is one of the most mature subfields of CS research. However, the more well-known of such algorithms have been designed with the expectation that the objects they sort can be compared in constant time. When used to sort objects that require linear-time comparison operations, such as strings, they perform a lot of wasterul work that leads to suboptimal performance. For maximum We have implemented a family of three different string sorting algorithms in Python and compared their performance against Python's native Timsort using a variety of different datasets.

Efficiencing string starting salgorithms are energed ded. A 100MB and a 200MB sample of each dataset TEINS, DNA and ENGLISH datasets from the Pizza&Chili Corpus, in addition to a set of URLs from Ranjan Sinha's ref1 data ref2 for his original Burstsort paper.

was used. The ENGLISH datasets were not used as-is, but with each word split on its own line, in order to make the algorithms sort individual words and not entire lines. The statistics file documents some stringological properties of these

ref1 https://sites.google.com/site/ranjansinha/home ref2 http://www.cs.mu.oz.au/rsinha/resources/data/so

ALGORITHMS

MSD RADIX SORT

MSD (most significant digit) radix sort is a divideand-conquer algorithm that partitions the strings based on their character at a given position. The comparison position starts from 0 and increases with one at every recursion level. No position, then, is visited twice; and if the algorithm does not attempt to partition buckets of size 1 or consisting entirely of strings shorter than the recursion depth, each string is visited at most one more time than the length of its shortest distinguishing prefix. Thus, the partitioning takes at most O(L(R) + n) time, where L(R) is the sum of the LCP array.

However, efficient implementations require the buckets to be implemented as an array of linked lists in order to avoid the overhead of binary search tree insertions and lookups. This allows true constant time insertion to buckets, but wastes time and memory if the strings use only a fraction of the alphabet for which MSD radix sort allocates space. Likewise, if the number of strings is smaller than the size of the alphabet, standard comparison based string sorting algorithms outperform MSD radix sort.

Our implementation uses a fixed alphabet size of 256 and falls back to ternary quicksort when the size of the bucket drops below it.

QUICKSORT ALGORITHMS

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BURST SORT

Burst sort text Burst sort text

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