Quantum algorithms for multivariate mean estimation

A. J. Cornelissen¹ Joint work with S. Jerbi² & Y. Hamoudi³

¹QuSoft, University of Amsterdam ²Institute for Theoretical Physics, University of Innsbruck ³Department of Electrical Engineering and Computer Sciences, University of California, Berkeley

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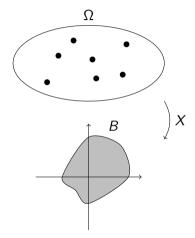




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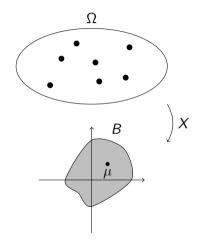
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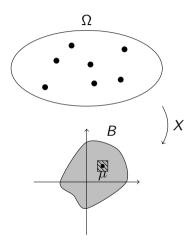


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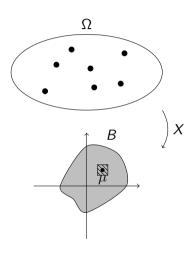
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Variants:

	Classically	Quantumly
Univariate $(d=1)$	Textbook	Textbook
Multivariate $(d>1)$	Textbook	Topic of this talk



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Calls to these routines are *samples*.

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Amplitude estimation finds $\mathbb{E}[X] \pm \varepsilon$ with $N = \mathcal{O}(1/\varepsilon)$ calls to U.



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Quadratic quantum speed-up!



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- Run 1D-algorithm with N samples:
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Roadblock: It is not clear how to represent $\mathbb{E}[X] \in \mathbb{R}^d$ as an amplitude.



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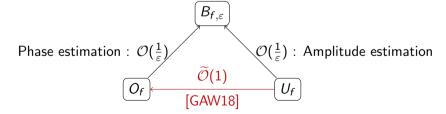
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- **1 Phase oracle:** $O_f: |x\rangle \mapsto e^{if(x)} |x\rangle$.

Intermezzo: Quantum oracle conversions [GAW18]

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Oracle conversion graph:





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0 1

Oan be turned into phase oracle:

$$O_f: |\mathbf{g}\rangle \mapsto e^{i\mathbf{g}\mathbb{E}[X]} |\mathbf{g}\rangle$$
 with $\widetilde{O}(1)$ calls to U_f .

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- **5** Recall quantum Fourier transform, for $k \in \{0, ..., 2^n 1\}$:

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- Oan be turned into phase oracle:
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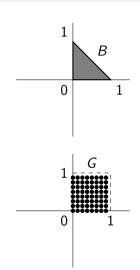
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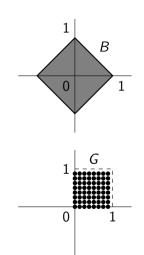


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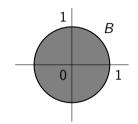
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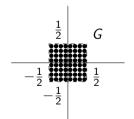


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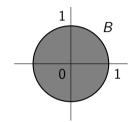


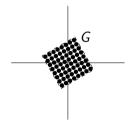


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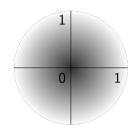


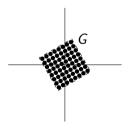


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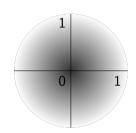


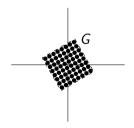
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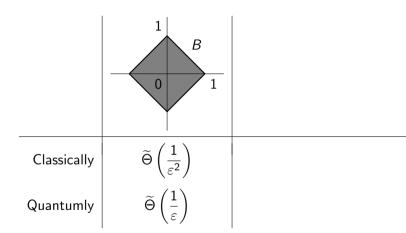
All improvements retain $\widetilde{O}(1/\varepsilon)$.



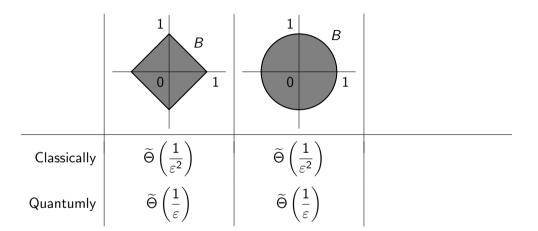


Limits

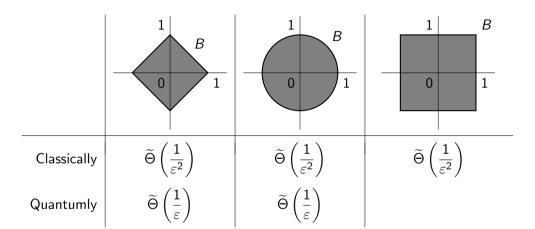
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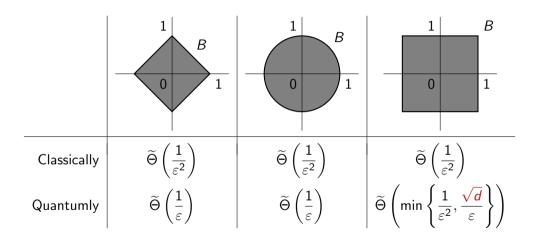
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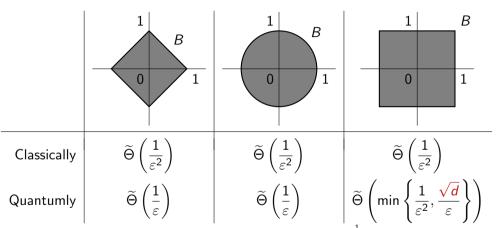


Limits



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Limits



For ℓ_p -approximations: multiply by $d^{\frac{1}{p}}$.

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Thanks for your attention! arjan@cwi.nl

