

Quantum approximate counting

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QuSoft, University of Amsterdam

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Quantum computing

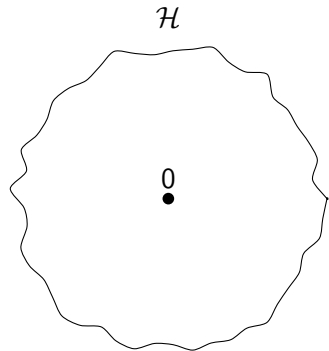
Quantum computing

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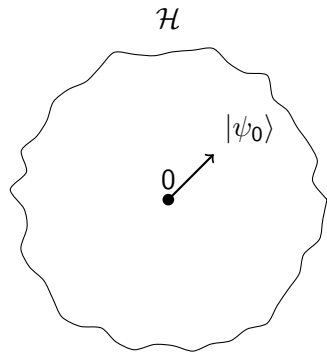
- 1 *State space*: Hilbert space \mathcal{H} .



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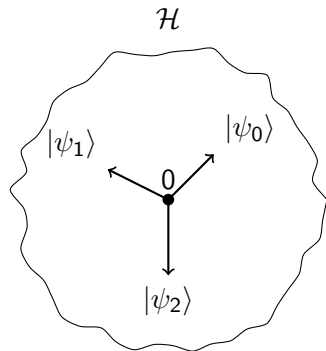
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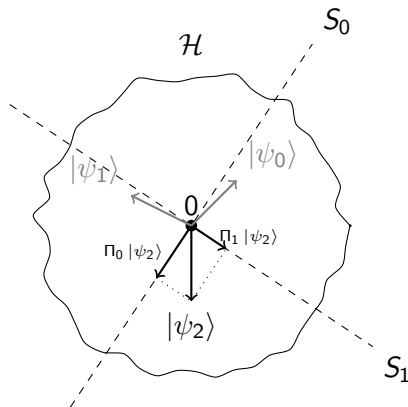


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$$\mathbb{P}(o) = \|\Pi_o |\psi_T\rangle\|^2.$$



Quantum computing

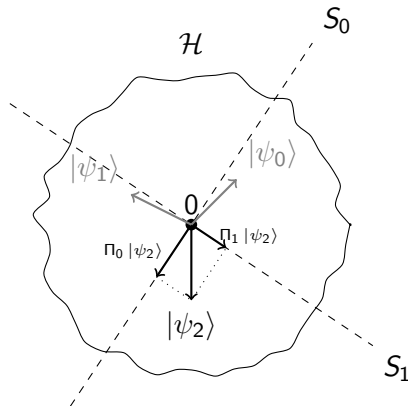
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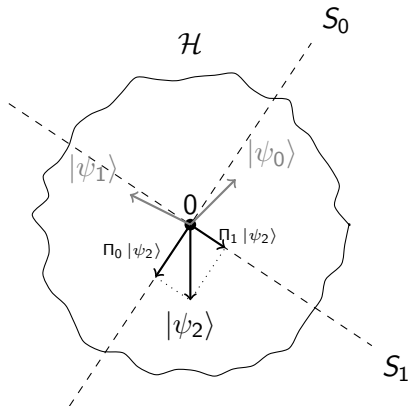
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$$|\psi_0\rangle \xrightarrow{U_1} |\psi_1\rangle \xrightarrow{O} |\psi_2\rangle \xrightarrow{U_3} |\psi_3\rangle \xrightarrow{O} \dots \xrightarrow{U_T} |\psi_T\rangle.$$



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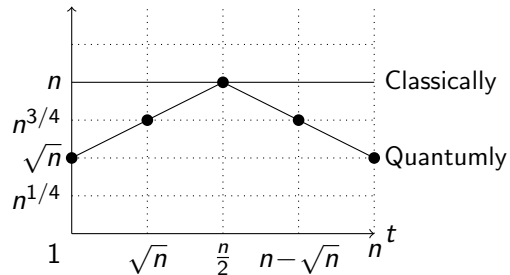
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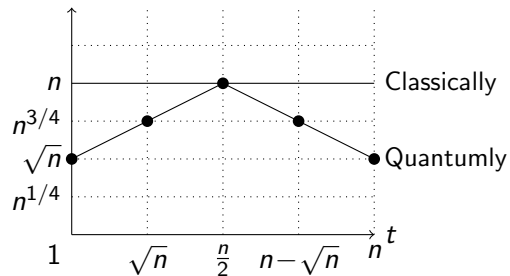
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- ④ *Goal for today*: look at the mathematics behind this phenomenon.

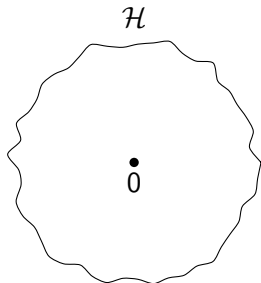
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Jordan's lemma

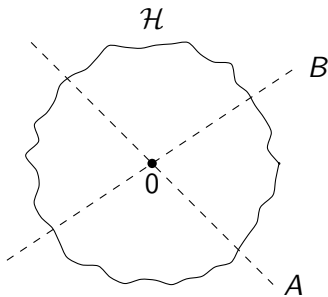
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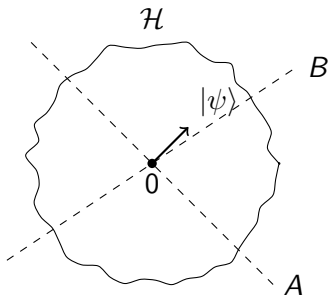
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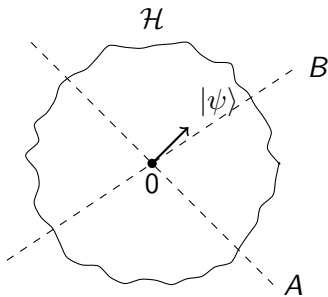
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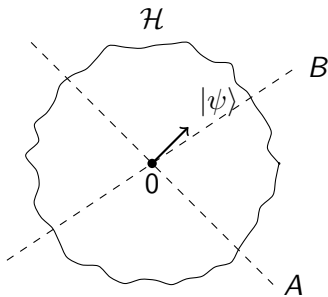
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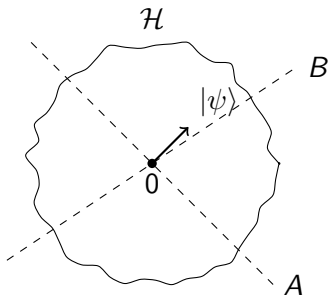


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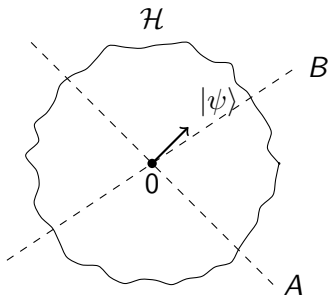


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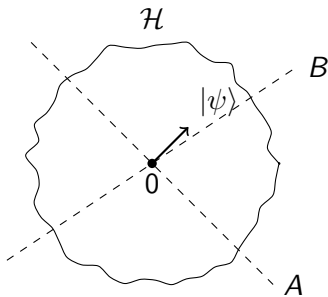


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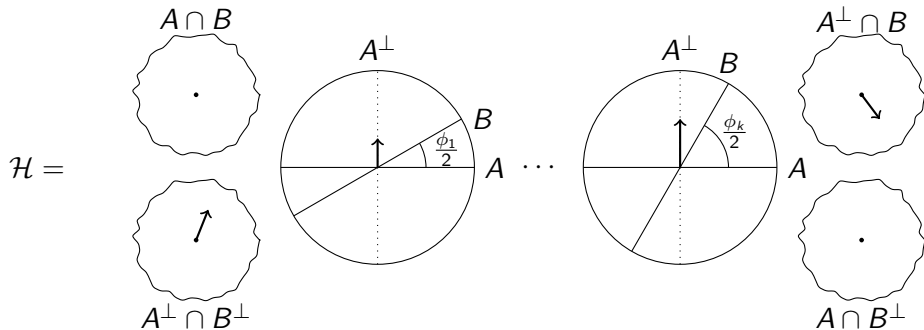


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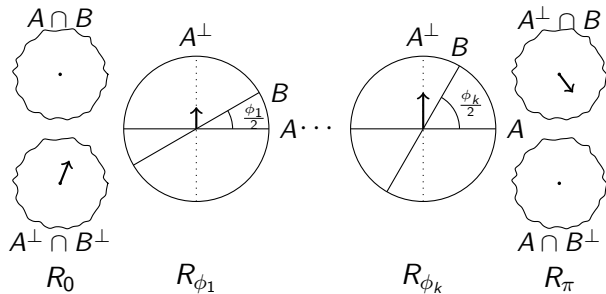
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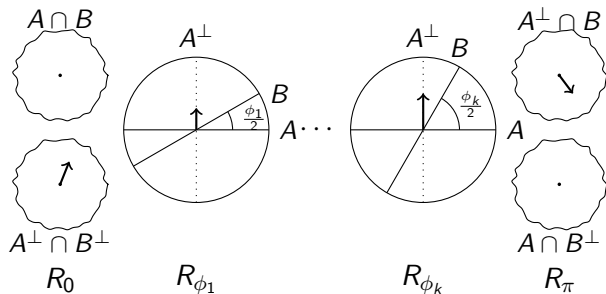
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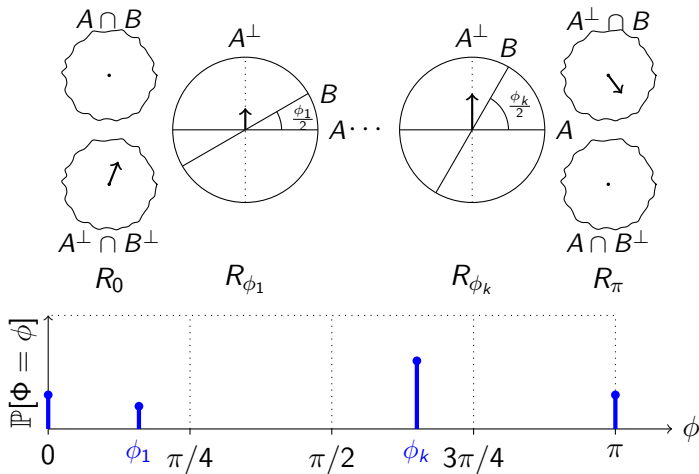
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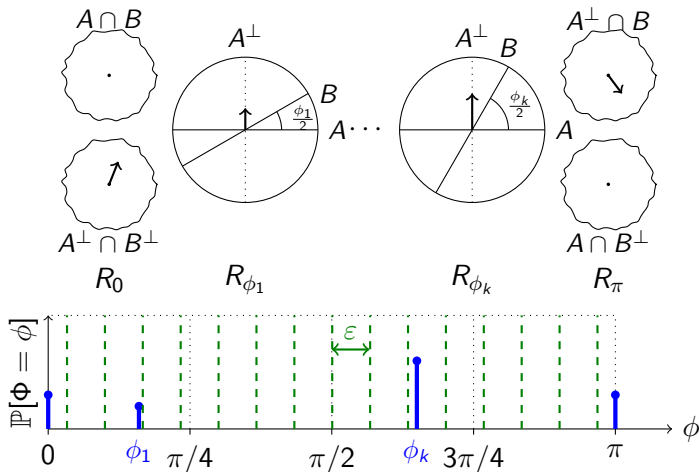
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- 3 **Phase estimation:**
One can sample from this binned distribution with $\mathcal{O}(1/\varepsilon)$ calls to $U = (2\Pi_B - I)(2\Pi_A - I)$.



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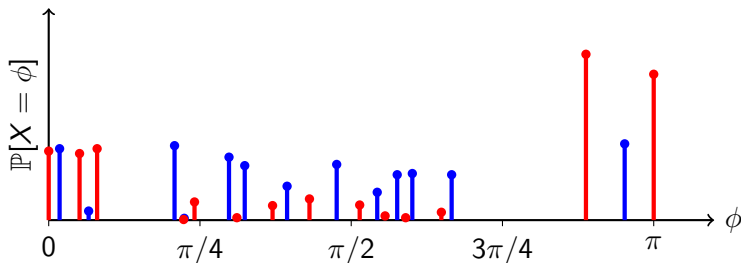
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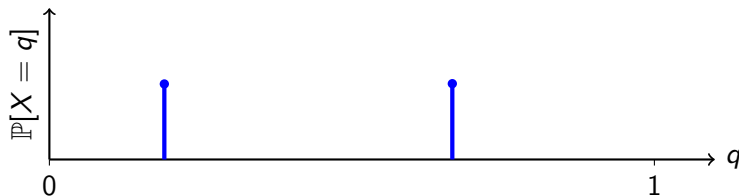
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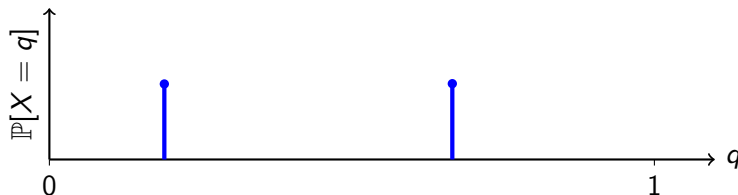
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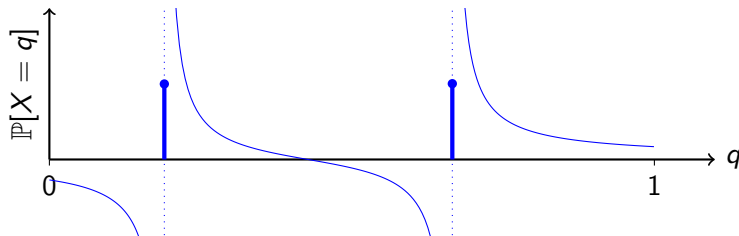
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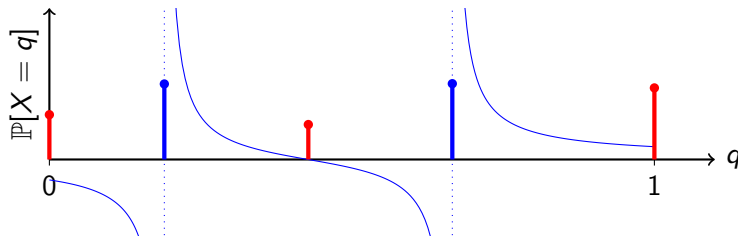
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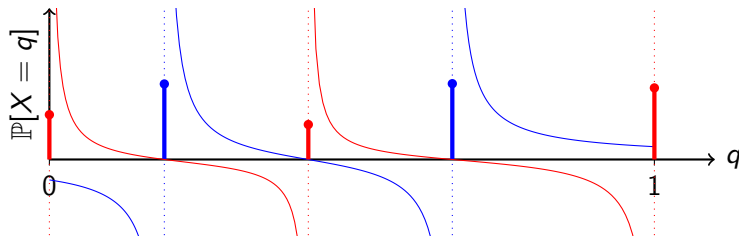
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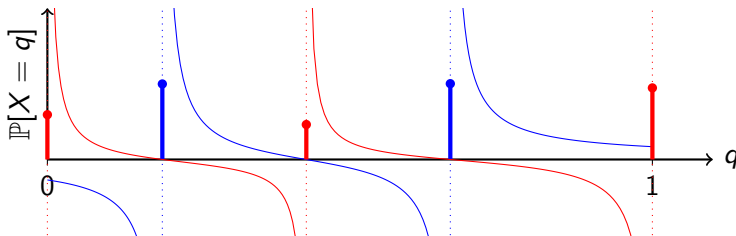
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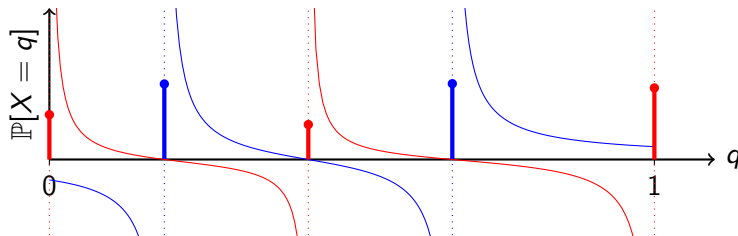
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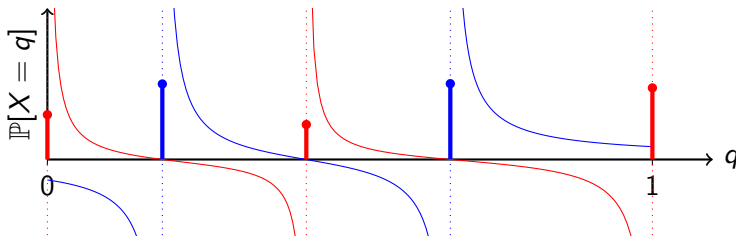
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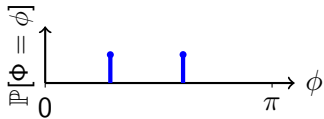
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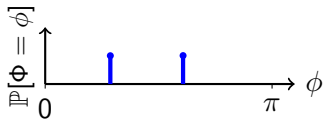
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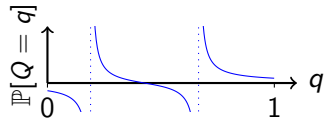
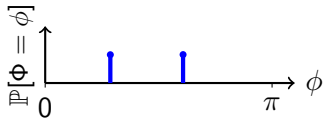
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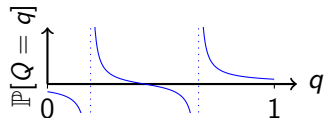
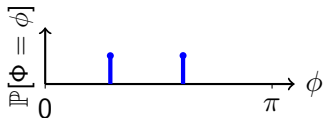
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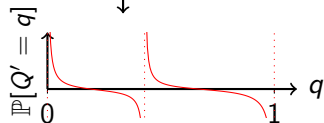
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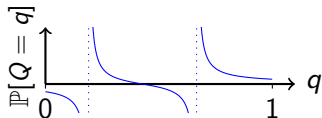
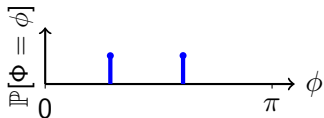
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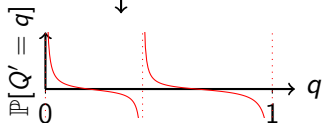
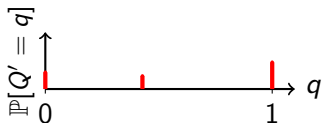
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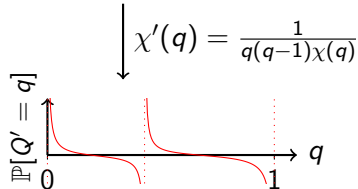
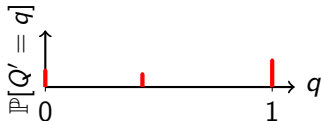
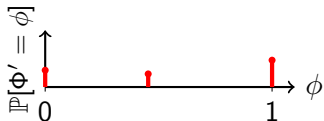
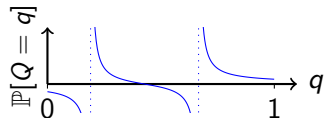
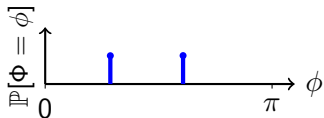


$$\mathbb{P}[Q' = q^*] = \lim_{q \rightarrow q^*} \chi'(q)(q - q^*)$$

Recap

$$Q = \sin^2(\Phi/2)$$

$$\chi(q) = \sum_{j=1}^k \frac{p_j}{q - q_j}$$



$$\Phi' = 2 \arcsin \sqrt{Q'}$$

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- 1 $f : \{1, 2, \dots, n\} \rightarrow \{0, 1\}.$
- 2 $O_f : |j\rangle \mapsto (-1)^{f(j)} |j\rangle.$
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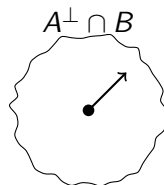
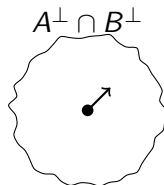
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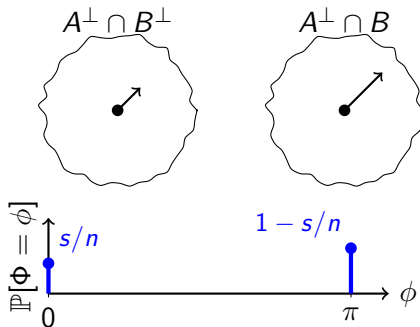
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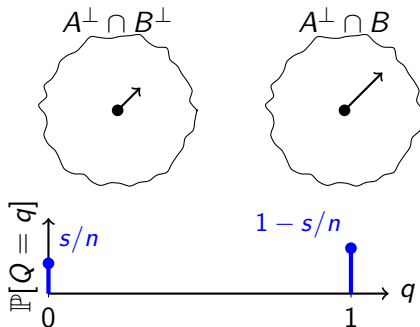
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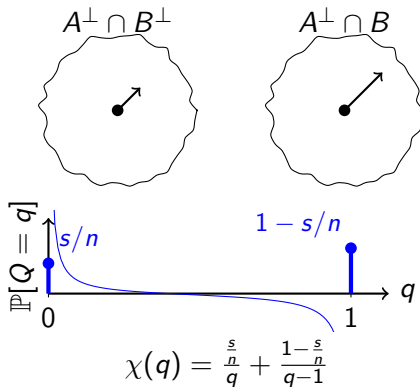
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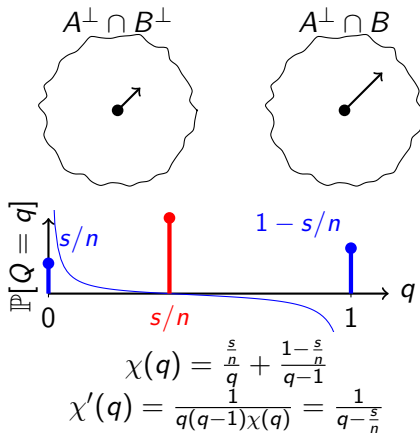
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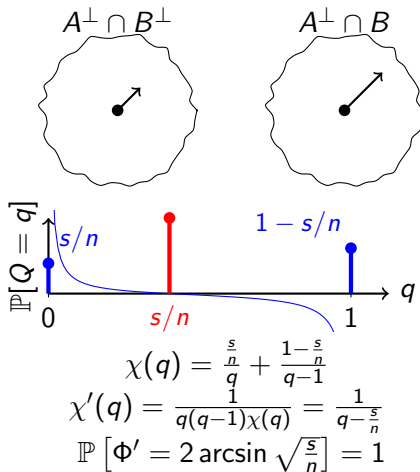
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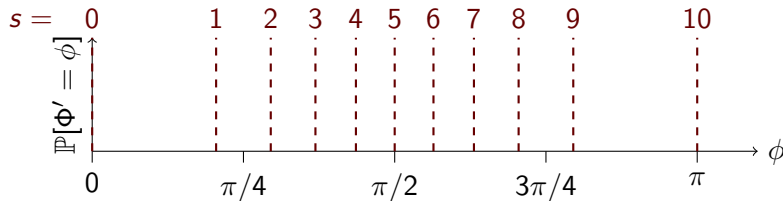


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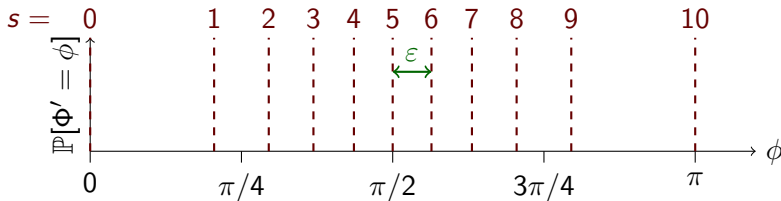


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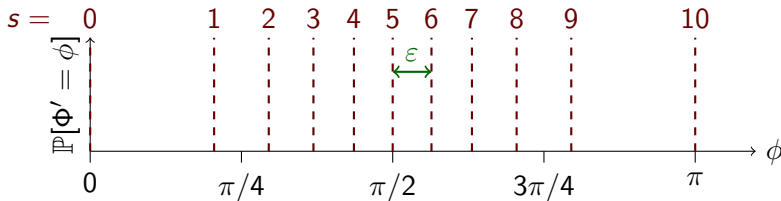


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Thus,

$$\varepsilon = \arcsin\left(\sqrt{\frac{t}{n}}\right) - \arcsin\left(\sqrt{\frac{t-1}{n}}\right) \Leftrightarrow \frac{1}{\varepsilon} = \mathcal{O}\left(\sqrt{t(n-t+1)}\right).$$

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Thanks for your attention!
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