

# Convex Optimization II

## Lecture 6: Convex Optimization Problems and Lagrange Duality

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# REFERENCE

[1] S. Boyd and L. Vandenberghe, *Convex Optimization*, first edition, Cambridge University Press, 2004.

All materials and figures in this lecture are from [1].

Thanks to Prof. Vincent Wong and Prof. Stephen Boyd for all the slides used in this lecture.

# OPTIMIZATION PROBLEM IN STANDARD FORM

$$\begin{aligned} & \text{minimize} && f_0(\mathbf{x}) \\ & \text{subject to} && f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & && h_i(\mathbf{x}) = 0, \quad i = 1, \dots, p \end{aligned} \tag{1}$$

- $\mathbf{x} \in \mathbf{R}^n$  is the optimization variable
- $f_0 : \mathbf{R}^n \rightarrow \mathbf{R}$  is the **objective** or **cost** function
- $f_i : \mathbf{R}^n \rightarrow \mathbf{R}, i = 1, \dots, m$ , are the **inequality constraint** functions
- $h_i : \mathbf{R}^n \rightarrow \mathbf{R}, i = 1, \dots, p$ , are the **equality constraint** functions
- Domain  $\mathcal{D} = \bigcap_{i=1}^m \text{dom } f_i \cap \bigcap_{i=1}^p \text{dom } h_i$

**Optimal value:**

$$p^* = \inf \{ f_0(\mathbf{x}) \mid f_i(\mathbf{x}) \leq 0, i = 1, \dots, m, \quad h_i(\mathbf{x}) = 0, i = 1, \dots, p \}$$

- $p^* = \infty$  if problem is infeasible (no  $\mathbf{x}$  satisfies the constraints)
- $p^* = -\infty$  if problem is unbounded below

# OPTIMAL AND LOCALLY OPTIMAL POINTS

- $\mathbf{x}$  is **feasible** if  $\mathbf{x} \in \mathcal{D}$  and it satisfies the constraints
- A feasible  $\mathbf{x}^*$  is **optimal** if  $f_0(\mathbf{x}^*) = p^*$ ;  $X_{\text{opt}}$  is the set of optimal points
- $\tilde{\mathbf{x}}$  is **locally optimal** if there is an  $R > 0$  such that  $\tilde{\mathbf{x}}$  is optimal for

$$\begin{array}{ll}\text{minimize} & f_0(\mathbf{x}) \\ \text{subject to} & f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m, \quad h_i(\mathbf{x}) = 0, \quad i = 1, \dots, p \\ & \|\mathbf{x} - \tilde{\mathbf{x}}\|_2 \leq R\end{array}$$

Examples: (with  $n = 1, m = p = 0$ )

- $f_0(x) = 1/x$ ,  $\text{dom } f_0 = \mathbf{R}_{++}$  :  $p^* = 0$ , but the optimal value is not achieved
- $f_0(x) = -\log x$ ,  $\text{dom } f_0 = \mathbf{R}_{++}$  :  $p^* = -\infty$ , this problem is unbounded below
- $f_0(x) = x \log x$ ,  $\text{dom } f_0 = \mathbf{R}_{++}$  :  $p^* = -1/e$ ,  $x^* = 1/e$
- $f_0(x) = x^3 - 3x$ ,  $p^* = -\infty$ , local optimal at  $x = 1$

# IMPLICIT CONSTRAINTS

- The standard form optimization problem has an **implicit constraint**

$$\mathbf{x} \in \mathcal{D} = \bigcap_{i=0}^m \text{dom } f_i \cap \bigcap_{i=1}^p \text{dom } h_i$$

- The constraints  $f_i(\mathbf{x}) \leq 0$  and  $h_i(\mathbf{x}) = 0$  are the explicit constraints.
- A problem is **unconstrained** if it has no explicit constraints ( $m = p = 0$ ).

Example:

$$\text{minimize } f_0(\mathbf{x}) = - \sum_{i=1}^k \log(b_i - \mathbf{a}_i^T \mathbf{x})$$

is an unconstrained problem with implicit constraints  $\mathbf{a}_i^T \mathbf{x} < b_i$

# FEASIBILITY PROBLEM

$$\begin{array}{ll}\text{find} & \mathbf{x} \\ \text{subject to} & f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & h_i(\mathbf{x}) = 0, \quad i = 1, \dots, p\end{array}$$

can be considered a special case of the general problem with  $f_0(\mathbf{x}) = 0$ .

$$\begin{array}{ll}\text{minimize} & 0 \\ \text{subject to} & f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & h_i(\mathbf{x}) = 0, \quad i = 1, \dots, p\end{array}$$

- $p^* = 0$  if the feasible set is nonempty; any feasible  $\mathbf{x}$  is optimal.
- $p^* = \infty$  if the feasible set is empty.

# CONVEX OPTIMIZATION PROBLEM

- Standard form convex optimization problem

$$\begin{aligned} & \text{minimize} && f_0(\mathbf{x}) \\ & \text{subject to} && f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & && a_i^T \mathbf{x} = b_i, \quad i = 1, \dots, p, \end{aligned}$$

where  $f_0, f_1, \dots, f_m$  are convex; equality constraints are affine.

- Often written as

$$\begin{aligned} & \text{minimize} && f_0(\mathbf{x}) \\ & \text{subject to} && f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & && A\mathbf{x} = \mathbf{b}. \end{aligned}$$

- Important property: the feasible set of a convex problem is convex.
- If the objective function  $g_0(\mathbf{x})$  is concave, then maximizing  $g_0(\mathbf{x})$  is equivalent to minimizing  $-g_0(\mathbf{x})$ .

# LOCAL AND GLOBAL OPTIMA

Any locally optimal point of a convex problem is (globally) optimal.

Proof:

- Suppose  $x$  is locally optimal for a convex problem. That is,  $x$  is feasible and

$$f_0(x) = \inf\{f_0(z) \mid z \text{ feasible}, \|z - x\|_2 \leq R\} \quad (2)$$

for some  $R > 0$ .

- Suppose  $x$  is not globally optimal. There is a feasible  $y$  such that  $f_0(y) < f_0(x)$ . Clearly,  $\|y - x\|_2 > R$ .
- Consider the point  $w = (1 - \theta)x + \theta y$ , where  $\theta = \frac{R}{2\|y - x\|_2}$ .
- We have  $\|w - x\|_2 = R/2 < R$ . By convexity of the feasible set,  $w$  is feasible.
- By convexity of  $f_0$ , we have

$$f_0(w) \leq (1 - \theta)f_0(x) + \theta f_0(y) < f_0(x),$$

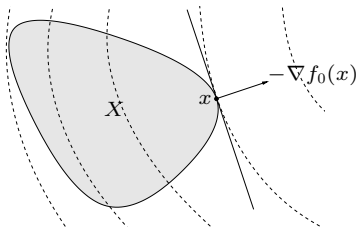
which contradicts (2). Hence, there exists no feasible  $y$  with  $f_0(y) < f_0(x)$ . That is,  $x$  is globally optimal.



# OPTIMALITY CRITERION FOR DIFFERENTIABLE $f_0$ (CONVEX PROBLEM)

$x$  is optimal if and only if it is feasible and

$$\nabla f_0(x)^T(y - x) \geq 0 \quad \text{for all feasible } y$$



if nonzero,  $\nabla f_0(x)$  defines a supporting hyperplane to feasible set  $X$  at  $x$

# SPECIAL CASES

- **unconstrained problem**

$$\nabla f_0(x) = 0$$

- **minimization over nonnegative orthant**

$$\text{minimize } f_0(x) \quad \text{subject to } x \succeq 0$$

The optimality condition can be expressed as

$$x \succeq 0, \quad \nabla f_0(x) \succeq 0, \quad x_i(\nabla f_0(x))_i = 0, \quad i = 1, \dots, n.$$

## EQUIVALENT PROBLEMS

We call two problems **equivalent** if from a solution of one, a solution of the other is readily found, and vice versa.

Example:

$$\begin{aligned} & \text{minimize} && \alpha_0 f_0(\mathbf{x}) \\ & \text{subject to} && \alpha_i f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & && \beta_i h_i(\mathbf{x}) = 0, \quad i = 1, \dots, p \end{aligned} \tag{3}$$

where  $\alpha_i > 0, i = 0, 1, \dots, m$ , and  $\beta_i \neq 0, i = 1, \dots, p$ .

Problems (1) and (3) are equivalent.

However, they are not the same unless  $\alpha_i$  and  $\beta_i$  are all equal to one.

General transformations that can yield equivalent problems include:

- Change of variables
- Transformation of objective and constraint functions
- Optimizing over some variables

## CHANGE OF VARIABLES

- Suppose  $\phi : \mathbf{R}^n \rightarrow \mathbf{R}^n$  is **one-to-one**, with image covering domain  $\mathcal{D}$  in problem (1).
- We define functions

$$\tilde{f}_i(\mathbf{z}) = f_i(\phi(\mathbf{z})), \quad i = 0, \dots, m, \quad \tilde{h}_i(\mathbf{z}) = h_i(\phi(\mathbf{z})), \quad i = 1, \dots, p.$$

- Consider the problem

$$\begin{aligned} & \text{minimize} && \tilde{f}_0(\mathbf{z}) \\ & \text{subject to} && \tilde{f}_i(\mathbf{z}) \leq 0, \quad i = 1, \dots, m \\ & && \tilde{h}_i(\mathbf{z}) = 0, \quad i = 1, \dots, p \end{aligned} \tag{4}$$

- Problems (1) and (4) are related by change of variable  $\mathbf{x} = \phi(\mathbf{z})$ .
- Problems (1) and (4) are equivalent.
- If  $\mathbf{x}$  solves problem (1), then  $\mathbf{z} = \phi^{-1}(\mathbf{x})$  solves problem (4).
- If  $\mathbf{z}$  solves problem (4), then  $\mathbf{x} = \phi(\mathbf{z})$  solves problem (1).

# TRANSFORMATION OF OBJECTIVE AND CONSTRAINT FUNCTIONS

- Suppose function  $\psi_0 : \mathbf{R} \rightarrow \mathbf{R}$  is **monotone increasing**. That is, for all  $x, y \in \text{dom } \psi_0$  such that  $x \leq y$ , we have  $\psi_0(x) \leq \psi_0(y)$
- Functions  $\psi_i, \dots, \psi_m : \mathbf{R} \rightarrow \mathbf{R}$  satisfy  $\psi_i(u) \leq 0$  if and only if  $u \leq 0$
- Functions  $\psi_{m+1}, \dots, \psi_{m+p} : \mathbf{R} \rightarrow \mathbf{R}$  satisfy  $\psi_i(u) = 0$  if and only if  $u = 0$
- We define functions

$$\tilde{f}_i(\mathbf{x}) = \psi_i(f_i(\mathbf{x})), \quad i = 0, \dots, m, \quad \tilde{h}_i(\mathbf{x}) = \psi_{m+i}(h_i(\mathbf{x})), \quad i = 1, \dots, p.$$

- Consider the problem

$$\begin{aligned} & \text{minimize} && \tilde{f}_0(\mathbf{x}) \\ & \text{subject to} && \tilde{f}_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & && \tilde{h}_i(\mathbf{x}) = 0, \quad i = 1, \dots, p \end{aligned} \tag{5}$$

- Problems (1) and (5) are equivalent.
- The feasible sets are identical. The optimal points are also identical.

## EXAMPLE: TRANSFORMATION OF OBJECTIVE FUNCTION

- Consider the unconstrained Euclidean norm minimization problem,

$$\text{minimize } \|A\mathbf{x} - \mathbf{b}\|_2,$$

with variable  $\mathbf{x} \in \mathbf{R}^n$ .

- Since the norm is non-negative, we can as well solve the problem

$$\text{minimize } \|A\mathbf{x} - \mathbf{b}\|_2^2 = (A\mathbf{x} - \mathbf{b})^T (A\mathbf{x} - \mathbf{b}).$$

- These two problems are equivalent.
- Note that the objective function in the first problem is not differentiable at any  $\mathbf{x}$  with  $A\mathbf{x} - \mathbf{b} = 0$ .
- The objective function in the second problem is differentiable for all  $\mathbf{x}$ .

## EXAMPLE: TRANSFORMATION OF CONSTRAINT FUNCTIONS

$$\begin{array}{ll}\text{minimize} & f_0(x) = x_1^2 + x_2^2 \\ \text{subject to} & f_1(x) = x_1/(1 + x_2^2) \leq 0 \\ & h_1(x) = (x_1 + x_2)^2 = 0\end{array}$$

- $f_0$  is convex; feasible set  $\{(x_1, x_2) \mid x_1 = -x_2 \leq 0\}$  is convex
- not a convex problem (according to our definition):  $f_1$  is not convex,  $h_1$  is not affine
- equivalent (but not identical) to the convex problem

$$\begin{array}{ll}\text{minimize} & x_1^2 + x_2^2 \\ \text{subject to} & x_1 \leq 0 \\ & x_1 + x_2 = 0\end{array}$$

## EXAMPLE: GEOMETRIC PROGRAMMING PROBLEM

- GP in standard form is **not** a convex optimization problem.

$$\begin{aligned} \underset{\mathbf{x} \succ \mathbf{0}}{\text{minimize}} \quad & \sum_{k=1}^{K_0} d_{0k} x_1^{a_{0k}^{(1)}} x_2^{a_{0k}^{(2)}} \cdots x_n^{a_{0k}^{(n)}} \\ \text{subject to} \quad & \sum_{k=1}^{K_i} d_{ik} x_1^{a_{ik}^{(1)}} x_2^{a_{ik}^{(2)}} \cdots x_n^{a_{ik}^{(n)}} \leq 1, \quad i = 1, \dots, m, \\ & d_l x_1^{a_l^{(1)}} x_2^{a_l^{(2)}} \cdots x_n^{a_l^{(n)}} = 1, \quad l = 1, \dots, M. \end{aligned}$$

- The main idea to convert a GP into a convex problem is based on a **logarithmic** change of variables, and a logarithmic transformation of the objective and constraint functions.
- Consider a **logarithmic** change of all the variables and multiplicative constants:

$$\begin{aligned} y_i &= \log x_i, \quad (\text{so } x_i = e^{y_i}), \quad i = 1, \dots, n, \\ b_{ik} &= \log d_{ik}, \quad (\text{so } d_{ik} = e^{b_{ik}}), \quad k = 1, \dots, K_i, \quad i = 0, 1, \dots, m, \\ b_l &= \log d_l, \quad (\text{so } d_l = e^{b_l}), \quad l = 1, \dots, M. \end{aligned}$$



## EXAMPLE: GEOMETRIC PROGRAMMING PROBLEM

- After the **logarithmic change of variables**, the problem becomes ( Why?)

$$\begin{aligned} & \underset{\mathbf{y}}{\text{minimize}} && \sum_{k=1}^{K_0} \exp(\mathbf{a}_{0k}^T \mathbf{y} + b_{0k}) \\ & \text{subject to} && \sum_{k=1}^{K_i} \exp(\mathbf{a}_{ik}^T \mathbf{y} + b_{ik}) \leq 1, && i = 1, \dots, m, \\ & && \exp(\mathbf{a}_l^T \mathbf{y} + b_l) = 1, && l = 1, \dots, M, \end{aligned}$$

where

$$\mathbf{a}_{ik}^T = \begin{bmatrix} a_{ik}^{(1)} & a_{ik}^{(2)} & \cdots & a_{ik}^{(n)} \end{bmatrix},$$

$$\mathbf{a}_l^T = \begin{bmatrix} a_l^{(1)} & a_l^{(2)} & \cdots & a_l^{(n)} \end{bmatrix},$$

and variables  $\mathbf{y} = (y_1, y_2, \dots, y_n)$ .

## EXAMPLE: GEOMETRIC PROGRAMMING PROBLEM

- After a logarithmic transformation of the objective and constraint functions, we have the following problem

$$\begin{aligned} & \underset{\mathbf{y}}{\text{minimize}} && \log \sum_{k=1}^{K_0} \exp(\mathbf{a}_{0k}^T \mathbf{y} + b_{0k}) \\ & \text{subject to} && \log \sum_{k=1}^{K_i} \exp(\mathbf{a}_{ik}^T \mathbf{y} + b_{ik}) \leq 0, && i = 1, \dots, m, \\ & && \mathbf{a}_l^T \mathbf{y} + b_l = 0, && l = 1, \dots, M. \end{aligned}$$

- Is the above problem convex? Why?

# OPTIMIZING OVER SOME VARIABLES

- We always have

$$\inf_{x,y} f(x,y) = \inf_x \tilde{f}(x)$$

where  $\tilde{f}(x) = \inf_y f(x,y)$ .

- We can always minimize a function by first minimizing over some of the variables, and then minimizing over the remaining ones.

## OPTIMIZING OVER SOME VARIABLES: EXAMPLE

- Consider a problem with strictly convex quadratic objective (i.e.,  $P_{11}, P_{22} \succeq 0$ ), with some of the variables unconstrained

$$\begin{aligned} \text{minimize} \quad & \mathbf{x}_1^T P_{11} \mathbf{x}_1 + 2\mathbf{x}_1^T P_{12} \mathbf{x}_2 + \mathbf{x}_2^T P_{22} \mathbf{x}_2 \\ \text{subject to} \quad & f_i(\mathbf{x}_1) \leq 0, \quad i = 1, \dots, m \end{aligned}$$

- We can first analytically minimize over  $\mathbf{x}_2$ :

$$\inf_{\mathbf{x}_2} (\mathbf{x}_1^T P_{11} \mathbf{x}_1 + 2\mathbf{x}_1^T P_{12} \mathbf{x}_2 + \mathbf{x}_2^T P_{22} \mathbf{x}_2) = \mathbf{x}_1^T (P_{11} - P_{12} P_{22}^{-1} P_{12}^T) \mathbf{x}_1$$

- The original problem is equivalent to

$$\begin{aligned} \text{minimize} \quad & \mathbf{x}_1^T (P_{11} - P_{12} P_{22}^{-1} P_{12}^T) \mathbf{x}_1 \\ \text{subject to} \quad & f_i(\mathbf{x}_1) \leq 0, \quad i = 1, \dots, m. \end{aligned}$$

# DUALITY MENTALITY

- Bound or solve an optimization problem via a different optimization problem.

standard form problem (not necessarily convex)

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p\end{array}$$

variable  $x \in \mathbf{R}^n$ , domain  $\mathcal{D}$ , optimal value  $p^\star$

**Lagrangian:**  $L : \mathbf{R}^n \times \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$ , with  $\text{dom } L = \mathcal{D} \times \mathbf{R}^m \times \mathbf{R}^p$ ,

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

- weighted sum of objective and constraint functions
- $\lambda_i$  is Lagrange multiplier associated with  $f_i(x) \leq 0$
- $\nu_i$  is Lagrange multiplier associated with  $h_i(x) = 0$

# LAGRANGE DUAL FUNCTION

- **Lagrange dual function:**  $g : \mathbf{R}^m \times \mathbf{R}^p \rightarrow \mathbf{R}$ ,

$$\begin{aligned} g(\lambda, \nu) &= \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) \\ &= \inf_{x \in \mathcal{D}} \left( f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right) \end{aligned}$$

$g$  is concave, can be  $-\infty$  for some  $\lambda, \nu$ .

- **Lower bound property:** if  $\lambda \succeq 0$ , then  $g(\lambda, \nu) \leq p^*$

Proof: if  $\tilde{x}$  is feasible and  $\lambda \succeq 0$ , we have

$$L(\tilde{x}, \lambda, \nu) = f_0(\tilde{x}) + \sum_{i=1}^m \lambda_i f_i(\tilde{x}) + \sum_{i=1}^p \nu_i h_i(\tilde{x}) \leq f_0(\tilde{x})$$

Thus,

$$g(\lambda, \nu) = \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) \leq L(\tilde{x}, \lambda, \nu) \leq f_0(\tilde{x})$$

Minimizing over all feasible  $\tilde{x}$  gives  $g(\lambda, \nu) \leq p^*$ .

# THE DUAL PROBLEM

## Lagrange dual problem

$$\begin{array}{ll}\text{maximize} & g(\lambda, \nu) \\ \text{subject to} & \lambda \succeq 0\end{array}$$

- finds best lower bound on  $p^*$ , obtained from Lagrange dual function
- a convex optimization problem; optimal value denoted  $d^*$
- $\lambda, \nu$  are dual feasible if  $\lambda \succeq 0, (\lambda, \nu) \in \mathbf{dom} g$
- often simplified by making implicit constraint  $(\lambda, \nu) \in \mathbf{dom} g$  explicit

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \succeq 0\end{array}$$

$$\begin{array}{ll}\text{maximize} & -b^T \nu \\ \text{subject to} & A^T \nu + c \succeq 0\end{array}$$



# WEAK AND STRONG DUALITY

**Weak Duality:**  $d^* \leq p^*$

- always holds for convex and nonconvex problems
- can be used to find nontrivial lower bounds for difficult problems

**Strong Duality:**  $d^* = p^*$

- does not hold in general
- (usually) holds for convex problems
- conditions that guarantee strong duality in convex problems are called **constraint qualifications**

**Duality Gap:**  $p^* - d^*$

- gives the gap between the optimal value of the primal problem (i.e., problem (1)) and the best (i.e., greatest) lower bound obtained from dual problem
- always non-negative

# SLATER'S CONSTRAINT QUALIFICATION

- strong duality holds for a convex problem

$$\begin{array}{ll}\text{minimize} & f_0(\mathbf{x}) \\ \text{subject to} & f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, m \\ & A\mathbf{x} = \mathbf{b}.\end{array}$$

- **constraint qualifications** provide conditions on the problem, beyond convexity, under which strong duality holds.
- **Slater's condition** : There exist  $\mathbf{x} \in \text{int } \mathcal{D}$  such that

$$f_i(\mathbf{x}) < 0, \quad i = 1, \dots, m, \quad A\mathbf{x} = \mathbf{b}.$$

- can be sharpened. If the first  $k$  constraint functions are affine, then strong duality holds provided the following condition holds: There exist  $\mathbf{x} \in \text{relint } \mathcal{D}$  (interior relative to affine hull) such that

$$f_i(\mathbf{x}) \leq 0, \quad i = 1, \dots, k, \quad f_i(\mathbf{x}) < 0, \quad i = k + 1, \dots, m, \quad A\mathbf{x} = \mathbf{b}.$$

# A NONCONVEX PROBLEM WITH STRONG DUALITY

- Strong duality holds for any optimization problem with quadratic objective and one quadratic inequality constraint, provided Slater's condition holds
- Example:

$$\begin{aligned} & \text{minimize} && x^T A x + 2b^T x \\ & \text{subject to} && x^T x \leq 1, \end{aligned}$$

where  $A \in \mathbf{S}^n$  and  $b \in \mathbf{R}^n$ .

When  $A \not\preceq 0$ , this is not a convex problem.

However, it can be shown that we have zero duality gap for this problem.

# COMPLEMENTARY SLACKNESS

assume strong duality holds,  $x^*$  is primal optimal,  $(\lambda^*, \nu^*)$  is dual optimal

$$\begin{aligned} f_0(x^*) = g(\lambda^*, \nu^*) &= \inf_x \left( f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^p \nu_i^* h_i(x) \right) \\ &\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*) \\ &\leq f_0(x^*) \end{aligned}$$

hence, the two inequalities hold with equality

- $x^*$  minimizes  $L(x, \lambda^*, \nu^*)$
- $\lambda_i^* f_i(x^*) = 0$  for  $i = 1, \dots, m$  (known as complementary slackness):

$$\lambda_i^* > 0 \implies f_i(x^*) = 0, \quad f_i(x^*) < 0 \implies \lambda_i^* = 0$$

# KARUSH-KUHN-TUCKER (KKT) CONDITIONS FOR NON-CONVEX PROBLEM

- Assume functions  $f_0, \dots, f_m, h_1, \dots, h_p$  are differentiable. They **may not be convex**.
- If strong duality holds and  $x^*, \lambda^*, \nu^*$  are optimal points, then they must satisfy the following **KKT** conditions:

$$f_i(x^*) \leq 0, \quad i = 1, \dots, m$$

$$h_i(x^*) = 0, \quad i = 1, \dots, p$$

$$\lambda_i^* \geq 0, \quad i = 1, \dots, m$$

$$\lambda_i^* f_i(x^*) = 0, \quad i = 1, \dots, m$$

$$\nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0.$$

- The last condition is due to the fact that  $x^*$  minimizes  $L(x, \lambda^*, \nu^*)$  over  $x$ , its gradient must vanish.

# KKT CONDITIONS FOR CONVEX PROBLEM

- If  $f_i$  are convex,  $h_i$  are affine,  $\tilde{x}, \tilde{\lambda}, \tilde{\nu}$  are points that satisfy the KKT conditions,

$$\begin{aligned}f_i(\tilde{x}) &\leq 0, & i = 1, \dots, m \\h_i(\tilde{x}) &= 0, & i = 1, \dots, p \\\tilde{\lambda}_i &\geq 0, & i = 1, \dots, m \\\tilde{\lambda}_i f_i(\tilde{x}) &= 0, & i = 1, \dots, m\end{aligned}$$

$$\nabla f_0(\tilde{x}) + \sum_{i=1}^m \tilde{\lambda}_i \nabla f_i(\tilde{x}) + \sum_{i=1}^p \tilde{\nu}_i \nabla h_i(\tilde{x}) = 0,$$

then  $\tilde{x}$  and  $(\tilde{\lambda}, \tilde{\nu})$  are primal and dual optimal, with zero duality gap.

- Since  $\tilde{\lambda}_i \geq 0$  and  $L(x, \tilde{\lambda}, \tilde{\nu})$  is convex in  $x$ , the last condition implies  $\tilde{x}$  minimizes  $L(x, \tilde{\lambda}, \tilde{\nu})$ . We have

$$g(\tilde{\lambda}, \tilde{\nu}) = L(\tilde{x}, \tilde{\lambda}, \tilde{\nu}) = f_0(\tilde{x}) + \sum_{i=1}^m \tilde{\lambda}_i f_i(\tilde{x}) + \sum_{i=1}^p \tilde{\nu}_i h_i(\tilde{x}) = f_0(\tilde{x})$$

- Thus,  $\tilde{x}$  and  $(\tilde{\lambda}, \tilde{\nu})$  have zero duality gap, and are primal and dual optimal.
- If a convex problem satisfies Slater's condition, then the KKT conditions provide **necessary and sufficient conditions** for optimality.

## EXAMPLE: WATER-FILLING (WITH $\alpha_i > 0$ )

$$\begin{array}{ll}\text{minimize} & -\sum_{i=1}^n \log(x_i + \alpha_i) \\ \text{subject to} & x \succeq 0, \quad \mathbf{1}^T x = 1\end{array}$$

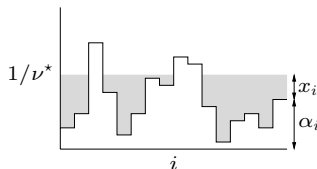
$x$  is optimal iff  $x \succeq 0$ ,  $\mathbf{1}^T x = 1$ , and there exist  $\lambda \in \mathbf{R}^n$ ,  $\nu \in \mathbf{R}$  such that

$$\lambda \succeq 0, \quad \lambda_i x_i = 0, \quad \frac{1}{x_i + \alpha_i} + \lambda_i = \nu$$

- if  $\nu < 1/\alpha_i$ :  $\lambda_i = 0$  and  $x_i = 1/\nu - \alpha_i$
- if  $\nu \geq 1/\alpha_i$ :  $\lambda_i = \nu - 1/\alpha_i$  and  $x_i = 0$
- determine  $\nu$  from  $\mathbf{1}^T x = \sum_{i=1}^n \max\{0, 1/\nu - \alpha_i\} = 1$

### interpretation

- $n$  patches; level of patch  $i$  is at height  $\alpha_i$
- flood area with unit amount of water
- resulting level is  $1/\nu^*$



# HOW TO SOLVE A CONVEX OPTIMIZATION PROBLEM?

- There are several efficient software to solve convex optimization problems
- Examples include CVX (<http://cvxr.com/cvx>), MOSEK (<http://www.mosek.com/>), CPLEX, and Matlab
- In MATLAB, you can use the optimization toolbox. The key function you need to use is *fmincon*:
- <https://www.mathworks.com/help/optim/ug/fmincon.html>
- $[x, fval, exitflag] = \text{fmincon}(\text{fun}, x0, A, b, Aeq, beq, lb, ub, \text{nonlcon}, \text{options})$

$$\begin{aligned} &\text{minimize} && f(x) \\ &\text{subject to} && c(x) \leq 0 \\ & && ceq(x) = 0 \\ & && A \cdot x \leq b \\ & && Aeq \cdot x = beq \\ & && lb \leq x \leq ub \end{aligned}$$



# SUMMARY

- **Convexity mentality**: Convex optimization is *nice* for several reasons:
  - ▶ local optimum is global optimum
  - ▶ zero duality gap (under technical conditions)
  - ▶ KKT optimality conditions are necessary and sufficient
- **Duality mentality**: Can always bound primal through dual, sometimes solve primal through dual
- Reading: Sections 4.2, 5.1, 5.2 and 5.5 in Boyd and Vandenberghe.

## EXAMPLE: GEOMETRIC PROGRAMMING PROBLEM

$$\begin{array}{ll}\underset{x_1, x_2}{\text{maximize}} & \log x_1 \log x_2 + \log x_1 \\ \text{subject to} & \log x_1 + e^{\log x_1 \log x_2} \leq 1\end{array}$$

$$\begin{array}{ll}\underset{x_1, x_2}{\text{minimize}} & -x_1 + x_2 \log x_2 \\ \text{subject to} & x_1 + x_2 = 1 \\ & x_1, x_2 \geq 0\end{array}$$