

# ASSIGNMENT 3

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Design and Analysis of Algorithm.

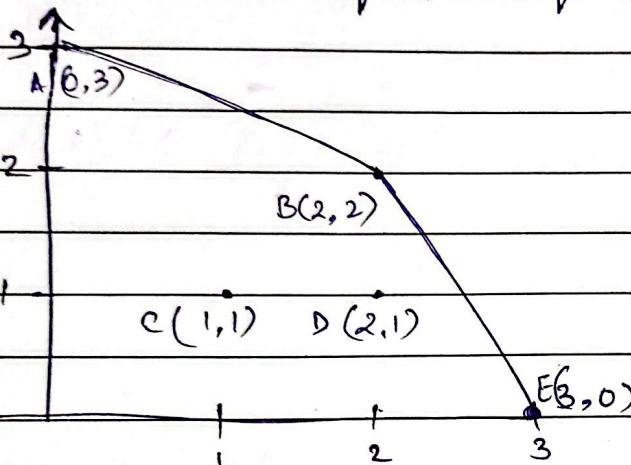
23K-0078

BAI-SA

Q1: A(0,3) B(2,2) C(1,1) D(2,1) E(3,0) - F(0,0)

Starting point: F(0,0)

Orientation:  $(P, q, r) = (q_x - p_x)(r_y - p_y) - (q_y - p_y)(r_x - p_x)$



Step	Current Point (P)	Candidate (Q)	Checking Point (R)	Orientation(P, Q, R)	Moving CW	New Candidate
1	F(0,0)	A(0,3)	B(2,2)	(0-0)(2-0)-3(0)(2)-6	CW	A
			C(1,1)	(0-0)(1-0)-3(0)(1)-3	CW	A
			D(2,1)	(0-0)(1-0)-3(0)(2)-6	CW	A
			E(3,0)	(0-0)(0-0)-3(0)(3)-9	CW	A
2	A(0,3)	B(2,2)	C(1,1)	(2-0)(1-3)-(2-3)(1-0)-3	CW	B
			D(2,1)	(2-0)(1-3)-(2-3)(2-0)-2	CW	B
			E(3,0)	(2-0)(0-3)-(2-3)(3-0)-3	CW	B
			F(0,0)	(2-0)(0-3)-(2-3)(0-0)-6	CW	B

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3	$B(2,2)$	$A(0,3)$	$C(1,1)$	$(6-2)(1-2) - (3-2)(1-2) + 3$	CCW	C
			$D(1,1)$	$(1-2)(1-2) - (1-2)(2-2) + 1$	CCW	D
			$E(2,1)$	$(2-2)(6-2) - (1-2)(3-2) + 1$	CCW	E
4.	$E(3,0)$	$A(0,3)$	$B(2,2)$	$(0-3)(2-0) - (3-0)(2-3) - 3$	CW	A
			$C(1,1)$	$(0-3)(1-0) - (3-0)(1-3) + 3$	CCW	C
			$D(2,1)$	$(1-3)(1-0) - (1-0)(2-3) - 1$	CW	C
			$F(0,0)$	$(1-3)(0-0) - (1-0)(0-3) + 3$	CCW	F

Convex Hull:  $F(0,0) \rightarrow E(3,0) \rightarrow B(2,2) \rightarrow A(0,3) \rightarrow F(0,0)$

O2: Brute force Algorithm

$$m = \text{len}(P)$$

for  $i$  in  $0 \dots m-1$ :

$$f[i] = 0$$

for  $k$  in range( $j, 0, -1$ ):

if  $P[0:k] == P[j-k+1:j+1]$ :

$$f[i] = k$$

break

Time Complexity :  $O(n^3)$

2. KMP Optimized Algorithm:

$$m = \text{len}(P)$$

$$A = [0] * n$$

$$K = 0$$

for  $i$  in  $1 \dots m+1$

while  $K > 0$  and  $P[K] != P[i]$ :

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$$K = A[K-1]$$

if  $P[K] = P[i]$

$$K = K + 1$$

$$A[i] = K$$

while computing  $A[i]$ , KMP uses the previously computed  $A[i-1]$  to avoid rechecking prefixes from scratch.

Time Complexity :  $\delta(n)$

3. Computation of  $A[i]$  for "ababaca" using both approaches  
Brute force.

i	$P[0 \dots i]$	K values.	$A[i]$
0	a	No proper Prefix $\rightarrow A[0]=0$	0
1	ab	$K=1; "a"$ vs " $b$ " $\rightarrow$ no $\rightarrow A[1]=0$	0
2	aba	$K=2; "ab"$ vs " $ba$ " $\rightarrow$ no; $K=1; "a"$ vs " $a$ " $\rightarrow$ match	1
3	abab	$K=3; "aba"$ vs " $bab$ " $\rightarrow$ no; $K=2; "ab"$ vs " $ab$ " $\rightarrow$ match	2
4	ababa	$K=4; "abab"$ vs " $baba$ " $\rightarrow$ no; $K=3; "aba"$ vs " $aba$ " $\rightarrow$ match	3
5	ababac	$K=5; "ababa"$ vs " $babac$ " $\rightarrow$ no; $K=4; "bab"$ vs " $bab$ " $\rightarrow$ no; $K=3; "aba"$ vs " $ba$ " $\rightarrow$ no; $K=2; "ab"$ vs " $ac$ " $\rightarrow$ no; $K=1; "a"$ vs " $c$ " $\rightarrow$ no	0
6	ababaca	$K=6; "ababac"$ vs " $babaca$ "; $K=5; "ababa"$ vs " $abaca$ " $\rightarrow$ no; $K=4; "bab"$ vs " $baca$ " $\rightarrow$ no; $K=3; "aba"$ vs " $aca$ " $\rightarrow$ no; $K=2; "ab"$ vs " $ca$ " $\rightarrow$ no; $K=1; "a"$ vs " $a$ " $\rightarrow$ match	1

$$F = \{0, 0, 1, 2, 3, 0, 1\}$$

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### Optimized KMP:

i	P[i]	K before	Compare ( $P[K] ? = P[i]$ )	Action	K after	A[i]
0	a	-	-	$F[0] = 0$	-	0
1	b	$K=F[0]=0$	$K=0 \rightarrow 'a' ? = 'b'$ → no	$K=0$	0	0
2	a	$K=F[0]=0$	$P[0] ? = P[2]: 'a' ? = 'a'$ → Y	$K=1$	1	1
3	b	$K=F[2]=1$	$P[1] ? = P[3]: 'b' ? = 'b'$ → Y	$K=2$	2	2
4	a	$K=F[3]=2$	$P[2] ? = P[4]: 'a' ? = 'a'$ → Y	$K=3$	3	3
5	c	$K=F[4]=3$	$P[3] ? = P[5]: 'b' ? = 'c'$ → no → $K=F[2]=1$	$K=1; P[1], K=0; P[0] ? = 0$ $'b' ? = 'c'$ → $? = 'c'$ → no → $K=0$ no → $K=0$	0	0
6	a	$K=F[5]=0$	$P[0] ? = P[6]: 'a' ? = 'a'$ → yes	$K=1$	1	1

$$F = [0, 0, 1, 2, 3, 0, 1]$$

### 4. Comparison:

(i) Time Complexity:

Brute force:  $O(n^3)$

Optimized KMP:  $O(n)$

### (ii) Character Comparisons:

Brute force algorithm does many repeated comparisons of the same prefixes and suffixes from the beginning.

KMP reuses previous meet matches and ensures each character is compared only a limited number of times.

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Q3:

i	dp[i]	Coin = 1				Coin = 5				Coin = 6				Coin = 8				dp[13] = 8
		before	dp[i] (aft)															
0	1	1															1	
1	0	1															1	
2	0	1															1	
3	0	1															1	
4	0	1															1	
5	0	1															2	
6	0	1															3	
7	0	1															3	
8	0	1															4	
9	0	1															4	
10	0	1															5	
11	0	1															6	
12	0	1															7	
13	0	1															8	

(b) str1 = KITTEN , str2 = SITTING

if  $\text{str1}[i-1] == \text{str2}[j-1]$

$\text{dp}[i][j] = \text{dp}[i-1][j-1]$

else :

delete      insert      substitution.

$$\text{dp}[i][j] = 1 + \min(\text{dp}[i-1][j], \text{dp}[i][j-1], \text{dp}[i-1][j-1])$$

~~K / O / I / T / T / E / N~~      ~~S / I / T / T / I / N / G~~

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j 0 1 2 3 4 5 6 7  
 " S I T T I N G  
 :

0'	0	1	2	3	4	5	6	7	$dp[6][7] = 3$
1 K	1	1	2	3	4	5	6	7	Substitute K → S
2 T	2	2	1	2	3	4	5	6	Substitute T → I
3 T	3	3	2	1	2	3	4	5	Insert G at the end.
4 T	4	4	3	2	1	2	3	4	
5 E	5	5	4	3	2	2	3	4	
6 N	6	6	5	4	3	3	2	3	

c). length [] = {1, 2, 3, 4, 5, 6, 7, 8}, price = {1, 5, 8, 9, 10, 16, 18, 20}

Reqd length = 8

j	i: $P_i + \sigma[j-i]$	r[j]	cut[i]
1	1: $P_1 + \sigma_0 \Rightarrow 1+0 \Rightarrow 1$	1	1
2	1: 1+1 $\Rightarrow 2$ ; 2: 5+0 $\Rightarrow 5$	5	2
3	1: 1+5 $\Rightarrow 6$ ; 2: 5+1 $\Rightarrow 6$ ; 3: 8+0 $\Rightarrow 8$	8	3
4	1: 1+8 $\Rightarrow 9$ ; 2: 5+5 $\Rightarrow 10$ ; 3: 8+9 $\Rightarrow 9$ ; 4: 9+0 $\Rightarrow 9$	10	2
5	1: 1+10 $\Rightarrow 11$ ; 2: 5+8 $\Rightarrow 13$ ; 3: 8+5 $\Rightarrow 13$ ; 4: 9+1 $\Rightarrow 10$ ; 5: 10+0 $\Rightarrow 10$	13	2
6	1: 1+13 $\Rightarrow 14$ ; 2: 5+10 $\Rightarrow 15$ ; 3: 8+8 $\Rightarrow 16$ ; 4: 9+5 $\Rightarrow 14$ ; 5: 10+1 $\Rightarrow 11$ ; 6: 16+0 $\Rightarrow 16$	16	3
7	1: 1+16 $\Rightarrow 17$ ; 2: 5+13 $\Rightarrow 18$ ; 3: 8+10 $\Rightarrow 18$ ; 4: 9+8 $\Rightarrow 17$ ; 5: 10+5 $\Rightarrow 15$ ; 6: 16+1 $\Rightarrow 17$ ; 7: 18+0 $\Rightarrow 18$	18	2
8	1: 1+18 $\Rightarrow 19$ ; 2: 5+16 $\Rightarrow 21$ ; 3: 8+13 $\Rightarrow 21$ ; 4: 9+10 $\Rightarrow 19$ ; 5: 10+8 $\Rightarrow 18$ ; 6: 16+5 $\Rightarrow 21$ ; 7: 18+1 $\Rightarrow 19$ ; 8: 20+0 $\Rightarrow 20$	21	2

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PRODUCTS

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$$r = [0, 1, 5, 8, 10, 13, 16, 18, 21]$$

$$\text{cut} = [0, 1, 2, 3, 2, 2, 3, 2, 2]$$

$$j = 8 \rightarrow \text{cut}[8] = 2 ; 8 - 2 \rightarrow 6$$

$$j = 6 \rightarrow \text{cut}[6] = 3 ; 6 - 3 \rightarrow 3$$

$$j = 3 \rightarrow \text{cut}[3] = 3 ; 3 - 3 \rightarrow 0$$

$$\text{Total Value} = P_2 + P_3 + P_3 \Rightarrow 5 + 8 + 8 \Rightarrow 21$$

(d) word dict = {i, like, ice, cream, icecream, mobile, apple}

input = i like apple.

j	dp[ij]	Substring	Valid Split	Valid Segmentation
0	True	" "	base case	
1	True	i	" ; "	" i like apple "
2	False	il	-	
3	False	ili	-	
4	False	ilik	-	
5	True	ilike	" ; like "	
6	False	ilikea	-	
7	False	ilikeap	-	
8	False	ilikeapp	-	
9	False	ilikeappl	-	
10	True	ilikeapple	" ; like apple "	

Q4: Scenario : Meal Planning within Budget

Goal: Maximize total nutrition while staying within a \$15 weekly food budget

Weight: Cost(money)

Capacity 15 dollars

Value: Nutrition gained

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### Data Table

Meals	Cost (\$)	Nutrition (Value)
Salad	4	8
Chicken Rice	6	12
Pasta	5	10
Soup	3	6
Smoothie	2	5

dp [Meal] [budget]

Meal ↓ / Budget →	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Salad (4, 8)	0	0	0	0	8	8	8	8	8	8	8	8	8	8	8
Chicken Rice (6, 12)	0	0	0	0	8	8	12	12	12	12	20	20	20	20	20
Pasta (5, 10)	0	0	0	0	8	10	12	12	12	18	20	22	22	22	22
Soup (3, 6)	0	0	0	6	8	10	12	14	16	18	18	22	24	24	26
Smoothie (2, 5)	0	0	5	6	8	10	12	14	17	19	21	23	25	26	27

Final: dp [5][15] = 27

Max Value = 27

Best Meals : Chicken Rice + Pasta + soup + Smoothie.

Total Cost:  $6 + 5 + 3 + 2 = 15$

Nutrition: 27

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