These lecture notes include some material from Professors Bertossi, Kolaitis, Guagliardo, Vardi, Libkin, Barland, McMahan

Nested Queries

Lecture Handout

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Aggregate results in WHERE

The right way

Account

Number	Branch	CustID	Balance
111	London	1	1330.00
222	London	2	1756.00
333	Edinburgh	1	450.00

Accounts with a higher balance than the average of all accounts

	Number
Answer:	111
	222

Aggregate results in WHERE

The wrong way

Accounts with a higher balance than the average of all accounts

ERROR in PostgreSQL

Aggregate functions can only be used in **SELECT** and **HAVING**

Check in SQLite

Comparisons with subquery results

```
SELECT
...
FROM
...
WHERE
term op ( subquery );
Allowed as long as subquery returns a single value

SELECT
...
FROM
...
WHERE
(term<sub>1</sub>, ..., term<sub>n</sub>) op ( subquery );
Allowed as long as subquery returns a single row with n columns
```

The WHERE clause revisited

 $term := attribute \mid value$

comparison :=

- **▶** (term, ..., term) **op** (term, ..., term) with **op** ∈ {=, <>, <, > <=, >=}
- ► term IS [NOT] NULL
- ► (term, ..., term) **op ANY** (query)
- ► (term, ..., term) **op ALL** (query)
- ► (term, ..., term) [NOT] IN (query)
- ► EXISTS (query)

condition :=

- comparison
- ► condition1 **AND** condition2
- ► condition1 **OR** condition2
- ▶ **NOT** condition

Comparisons between tuples

$$(A_{1}, \dots A_{n}) = (B_{1}, \dots, B_{n}) \iff A_{1} = B_{1} \wedge \dots \wedge A_{n} = B_{n}$$

$$(A_{1}, \dots A_{n}) <> (B_{1}, \dots, B_{n}) \iff A_{1} \neq B_{1} \vee \dots \vee A_{n} \neq B_{n}$$

$$(A_{1}, A_{2}, A_{3}) < (B_{1}, B_{2}, B_{3}) \qquad \text{(generalizes to } n \text{ elements)}$$

$$\iff A_{1} < B_{1} \vee \left(A_{1} = B_{1} \wedge \left(A_{2} < B_{2} \vee \left(A_{2} = B_{2} \wedge A_{3} < B_{3}\right)\right)\right)$$

$$(A_{1}, A_{2}, A_{3}) <= (B_{1}, B_{2}, B_{3}) \qquad \text{(generalizes to } n \text{ elements)}$$

$$\iff A_{1} < B_{1} \vee \left(A_{1} = B_{1} \wedge \left(A_{2} < B_{2} \vee \left(A_{2} = B_{2} \wedge A_{3} \leq B_{3}\right)\right)\right)$$

```
(term, ..., term) op ANY ( query )
```

True if **there exists** a row \bar{r} in the results of query such that (term, ..., term) **op** \bar{r} is true

Examples:

- ▶ $3 < ANY(\{1, 2, 3\})$ is false
- ► $3 < ANY(\{2, 3, 4\})$ is true
- ▶ What about $3 < ANY({})?$

ALL

```
(term, ..., term) op ALL ( query )
```

True if for all rows \bar{r} in the results of query (term, ..., term) op \bar{r} is true

Examples:

- ▶ $3 < ALL(\{5,4,6\})$ is true
- ▶ $3 < ALL({4,3,5})$ is false
- ▶ What about $3 < \mathbf{ALL}(\{\})$?

Examples with ANY / ALL

ID of customers from London who own an account

Customers living in cities without a branch

IN / NOT IN

```
(term, ..., term) IN ( query )
    same as
(term, ..., term) = ANY ( query )

(term, ..., term) NOT IN ( query )
    same as
(term, ..., term) <> ALL ( query )
```

Examples with IN / NOT IN

ID of customers from London who own an account

Customers living in cities without a branch

EXISTS

```
EXISTS ( query ) is true if the result of query is non-empty
```

```
(Stupid) Example
```

Return all the customers if there are some accounts in London

Correlated subqueries

All nested queries can refer to attributes in the parent queries

```
(Smarter) Example
```

Return customers who have an account in London

parameters = attributes of a subquery that refer to outer queries

Examples with EXISTS / NOT EXISTS

ID of customers from London who own an account

Customers living in cities without a branch

Scoping

A subquery has

- ► a **local scope** (its **FROM** clause)
- ▶ n outer scopes (where n is the level of nesting) (these are the **FROM** clauses of the parent queries)

For each reference to an attribute

- 1. Look for a binding in the local scope
- 2. If no binding is found, look in the closest outer scope
- 3. If no binding is found, look in the next closest outer scope
- 4. ...
- 5. If no binding is found, give error

Attribute bindings

What A, B refer to depends on the attributes in table1 and table2

- Always give aliases to tables
- Always prefix the attributes with the tables they refer to

The FROM clause revisited

```
FROM table<sub>1</sub> [[AS] T_1], ..., table<sub>n</sub> [[AS] T_n]

table :=

base-table

join-table

(query)

join-table :=

table JOIN table ON condition

table NATURAL JOIN table

table CROSS JOIN table
```

Subqueries in FROM

```
Must always be given a name (in standard SQL)

SELECT * FROM ( SELECT * FROM R );

ERROR: subquery in FROM must have an alias

Cannot refer to attributes of other tables in the same FROM clause

SELECT *

FROM R, ( SELECT * FROM S WHERE S.a=R.a ) S1;

ERROR: invalid reference to FROM-clause entry for table "r"
```

Example: Queries with and without HAVING

Branches with a total balance (across accounts) of at least 500

```
SELECT
          A.branch
FROM
          Account A
GROUP BY A.branch
          SUM(A.balance) >= 500 ;
HAVING
Same query without HAVING:
SELECT subquery.branch
FROM
       ( SELECT
                  A.branch, SUM (A.balance) AS total
         FROM
                  Account A
         GROUP BY A.branch ) AS subquery
       subquery.total >= 500;
WHERE
```

Example: Aggregation on aggregates

Average of the total balances across each customer's accounts

- 1. Find the total balance across each customer's accounts
- 2. Take the average of the totals

Other bits and pieces of SQL

Ordering

```
ORDER BY \langle \text{column}_1 \rangle [DESC], ..., \langle \text{column}_n \rangle [DESC]

Sorts the output rows according to the values of \text{column}_1

If two rows have the same value for \text{column}_1, they are sorted by the values of \text{column}_2 and so on ...
```

- ► Default ordering is **ascending** (can be specified with **ASC**)
- ► **Descending** ordering is specified by **DESC**

Ordering example (1)

Account

Number	Branch	CustID	Balance
111	London	1	1330.00
222	London	2	1756.00
333	Edinburgh	1	450.00

SELECT 9

FROM Account

ORDER BY custid ASC, balance DESC;

Number	Branch	CustID	Balance
111	London	1	1330.00
333	Edinburgh	1	450.00
222	London	2	1756.00

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Ordering example (2)

Account

Number	Branch	CustID	Balance
111	London	1	1330.00
222	London	2	1756.00
333	Edinburgh	1	450.00

SELECT *

FROM Account

ORDER BY custid DESC, balance ASC;

Number	Branch	CustID	Balance
222	London	2	1756.00
333	Edinburgh	1	450.00
111	London	1	1330.00

Casting

```
In most dialects of SQL:
```

```
CAST( term AS \langle type \rangle )
```

Rounding

```
CAST (102, 4675 AS NUMERIC (5, 2)) gives 102.47 Useful also to produce values in a specific format
```

Aggregation

```
AVG( CAST( term AS NUMERIC(p,s) ) avoids rounding errors in some systems
```

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Acknowledgements

- [1] Database Systems: The Complete Book, 2nd EditionHector Garcia-Molina, Jeffrey D. Ullman, Jennifer WidomPrentice Hall, 2009
- [2] Database System Concepts, Seventh EditionAvi Silberschatz, Henry F. Korth, S. SudarshanMcGraw-Hill, March 2019www.db-book.com

Additional references and resources used in preparation of this course are listed on the course webpage or mentioned in slides.