

# CS251 - Data Structures and Algorithms

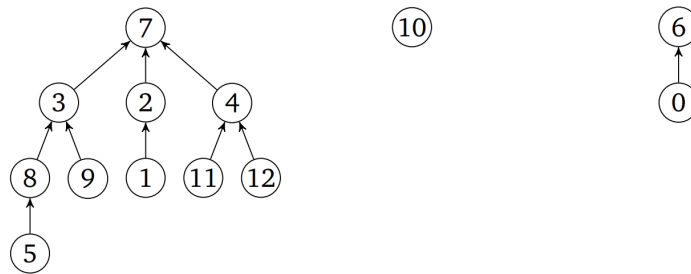
## Fall 2024

PSO 10, Week 12

### Question 1

#### (Union find)

Consider the following trees, which are a part of a disjoint set.



For the following problems, use both the **union-by-weight** and **path compression**.

- (1) Draw the resulting tree(s) after calling `find(5)`. What value does the method return?
- (2) Draw the resulting tree(s) of calling `union(2,6)` on the result of (1).

**Question 2****(Minimum spanning tree)**

You have found a MST  $T$  of a huge Graph  $G$ .  $G$  has an edge  $(v_1, v_2)$  whose weight is 20. After the MST was found the edge weight of  $(v_1, v_2)$  is updated. How would you update your MST in the following cases? What's the runtime of your solution? (You may assume that all edges in  $G$  have distinct weights before and after we update the weight of the edge  $(v_1, v_2)$ .)

**Note:** We are looking for solutions that will update  $T$  instead of building a new spanning tree of the updated graph.

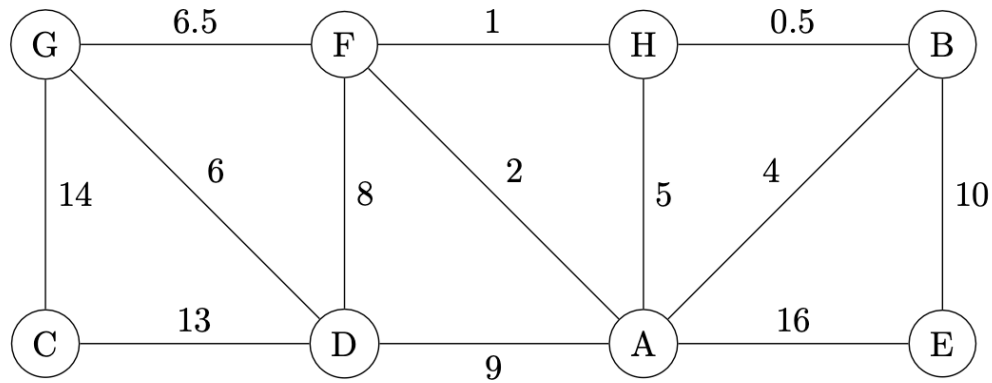
1.  $(v_1, v_2)$  was in  $T$  and the edge weight was updated to 5.
2.  $(v_1, v_2)$  was in  $T$  and the edge weight was updated to 40.
3.  $(v_1, v_2)$  was not in  $T$  and the edge weight was updated to 15.
4.  $(v_1, v_2)$  was not in  $T$  and the edge weight was updated to 25.

**Question 3****(More on the MST)**

1. An edge is called a **light-edge** crossing a cut  $\mathcal{C} := (S, V - S)$ , if its weight is the minimum of any edge crossing the cut. Show that:
  - If an edge  $(u, v)$  is contained in some MST (note that MST may not be unique), then it is a light-edge crossing some cut of the graph.
  - The converse is not true by giving a simple counter-example of a connected graph such that there exists a cut  $\mathcal{C} := (S, V - S)$ , in which  $(u, v)$  is a light-edge crossing the cut  $\mathcal{C}$  but does not form a MST of the graph.
2. Show that a graph has a unique MST, if for every cut of the graph, there is a unique light-edge crossing the cut. Show that the converse is not true by giving a counter-example.

**Question 4****(Kruskal's algorithm)**

1. Consider the following undirected graph. Assume that the graph is represented in adjacency-list form and that each adjacency-list is given in lexicographic order. List the order that edges are added when we run **Kruskal's algorithm**.



2. Suppose that all edge weights in a graph are integers in the range from 1 to  $|V|$ . How fast can you make Kruskal's algorithm run?