# CS251 - Data Structures and Algorithms Fall 2024

PSO 5, Week 6

```
1: function ExchangeSort(A : array)
      let n be the size of A
2:
      for i from 0 to n-2 do
3:
          for j from i+1 to n-1 do
4:
             if A[i] > A[j] then
5:
                SWAP(A, i, j)
6:
7:
             end if
          end for
8:
9:
      end for
      return A
10:
11: end function
```

```
1: function BubbleSort(A : array)
       let n be the size of A
3:
       repeat \leftarrow true
       while repeat do
4:
           repeat \leftarrow false
5:
           for i from 0 to n-2 do
6:
               if A[i] > A[i+1] then
7:
                  SWAP(A, i, i + 1)
9:
                  repeat \leftarrow true
               end if
10:
           end for
11:
12:
       end while
       return A
13:
14: end function
```

#### Question 1

The closed-form runtime expression T(n) for the number of compares between array items executed by ExchangeSort is:

A. 
$$T(n) = \frac{1}{2}n^2 - \frac{1}{2}n$$
  
B.  $T(n) = \frac{1}{2}n^2 + \frac{1}{2}n$   
C.  $T(n) = n^2 - 1$   
D.  $T(n) = n^2 + 1$   
E.  $T(n) = n^2$ 

# Question 2

The closed-form runtime expression T(n) for the maximum number of SWAP calls made by EXCHANGE-SORT is:

A. 
$$T(n) = \frac{1}{2}n^2 - \frac{1}{2}n$$
  
B.  $T(n) = \frac{1}{2}n^2 + \frac{1}{2}n$   
C.  $T(n) = n^2 - 1$   
D.  $T(n) = n^2 + 1$   
E.  $T(n) = n^2$ 

#### Question 3

The closed-form runtime expression T(n) for the maximum number of SWAP calls made by BUBBLESORT is:

A. 
$$T(n) = \frac{1}{2}n^2 - \frac{1}{2}n$$
  
B.  $T(n) = \frac{1}{2}n^2 + \frac{1}{2}n$ 

- C.  $T(n) = n^2 1$
- D.  $T(n) = n^2 + 1$
- E.  $T(n) = n^2$

# Question 4

If 
$$f(n) \in \Omega(g(n))$$
, then  $4^{f(n)} \in \Omega(4^{g(n)})$ .

- A. Necessarily True
- B. Possibly True
- C. Necessarily False

The Arrange function below rearranges the input array A content so that positive numbers are in even indices and negative numbers are in odd indices. The correctness of the algorithm requires that A be an array of size n, where n is an even integer. Also, A must contain n/2 positive integers and n/2 negative integers.

```
1: function Arrange(A : array)
        let n be the size of A
        for i from 0 to n-1 do
 3:
             if \neg((i \text{ is even } \land A[i] > 0) \lor (i \text{ is odd } \land A[i] < 0)) then
 4:
                 j \leftarrow i + 1
 5:
                 repeat \leftarrow true
 6:
                 while repeat \land j < n do
 7:
                     if (i is even \land A[j] > 0) \lor (i is odd \land A[j] < 0) then
 8:
9:
                         SWAP(A, i, j)
                         repeat \leftarrow false
10:
                     end if
11:
                     j \leftarrow j + 1
12:
13:
                 end while
             end if
14:
15:
         end for
         return A
16:
17: end function
```

#### Question 5

The following are arrays of size n = 6. Which input array will make Arrange to call SWAP n/2 times?

```
A. [1, -2, 3, 4, -5, -6]
B. [1, 2, 3, -4, -5, -6]
C. [-1, 2, -3, 4, -5, 6]
D. [-1, -2, 3, 4, 5, -6]
E. [-1, -2, 3, 4, -5, 6]
```

#### Question 6

Let K be an array of size n where the first n/2 items are negative, and the last n/2 items are positive. For instance, K = [-1, -2, -3, 4, 5, 6] with n = 6. The runtime expression T(n) for the number of times the condition in line 8 of the algorithm is evaluated when we call Arrange(K) is:

```
A. T(n) = \frac{1}{8}n^2 - \frac{1}{4}n

B. T(n) = \frac{1}{8}n^2 + \frac{1}{4}n

C. T(n) = \frac{1}{4}n^2 + \frac{1}{2}n

D. T(n) = \frac{1}{2}n^2 - \frac{1}{2}n

E. T(n) = \frac{1}{2}n - \frac{1}{2}
```

Consider a linked list where each node has an item and a pointer to the next node. A linked list has a cycle if at least one node exists in the list whose next pointer eventually points back to a previous node, forming a closed loop. The following is a cycle detection algorithm commonly known as the Tortoise and Hare algorithm:

```
1: function HasCycle(head : node)
         slow \leftarrow head
 2:
         fast \leftarrow head
 3:
         while fast \neq \text{null } \land fast.\text{next} \neq \text{null } \mathbf{do}
 4:
             slow \leftarrow slow.next
 5:
             fast \leftarrow fast.\text{next.next}
 6:
             if slow = fast then
 7:
                  return true
 8:
9:
             end if
         end while
10:
         return false
12: end function
```

#### Question 7

Suppose you run this function on a list with n nodes, with a cycle starting at node k (where  $1 \le k \le n$ ). What can be said about the number of steps the *slow* and *fast* pointers take before they meet?

- A. The slow and fast pointers will meet exactly after n steps.
- B. The slow pointer will always take n steps before meeting the fast pointer, regardless of where the cycle starts.
- C. The *slow* pointer will always enter the cycle before the *fast* pointer.
- D. The number of steps taken by the *slow* pointer before they meet depends on both the length of the list before the cycle and the length of the cycle itself.
- E. The slow and fast pointers may never meet if the cycle starts too close to the head of the list.

#### Question 8

Let Q in the following pseudocode be a fixed-capacity queue implemented with a circular array (i.e., the elements wrap around to reuse space at the front, as discussed in class).

```
1: function UNNECESSARYALGORITHM(n : \mathbb{Z}^+)
      Let Q be a queue implemented in an array of capacity n
2:
3:
      for i from 0 to 4 do
          Q.enqueue(i)
 4:
      end for
5:
      for i from 0 to 2 do
6:
          Q.dequeue()
 7:
       end for
8:
       for i from 5 to 11 do
9:
          Q.enqueue(i)
10:
       end for
11:
       Q.dequeue()
12:
       Q.enqueue(15)
13:
14: end function
```

At which index in the array is element 15 after calling UNNECESSARYALGORITHM(10)?

- A. 2
- B. 4
- C. 6
- D. 8
- E. 0

Let T be the binary tree for which the following traversals hold:

Preorder: T U T S R C E U R S
Inorder: S T R U C T U R E S
Postorder: S R T C U R U S E T

## Question 9

The height of T is:

- A. 2
- B. 3
- C. 4
- D. 5
- E. 6

## Question 10

T is:

- A. A balanced binary tree
- B. A complete binary tree
- C. A min binary heap
- D. A max binary heap
- E. A full binary tree

# Question 11

The most populated level of T is:

- A. 0
- B. 1
- C. 2
- D. 3
- E. 4

#### Question 12

Let B be the binary tree representation of a max binary heap. The values [8, 11, 3, 10, 7, 4, 6] are inserted in B (one at that time in the given order). What is the inorder traversal of B?

# Question 13

Let B' be the binary tree representation of a max binary heap built using HEAPIFY([8, 11, 3, 10, 7, 4, 6]). What is the inorder traversal of B'?