

# USENIX ATC 2023

## Towards Iterative Relational Algebra on the GPU

### Authors:

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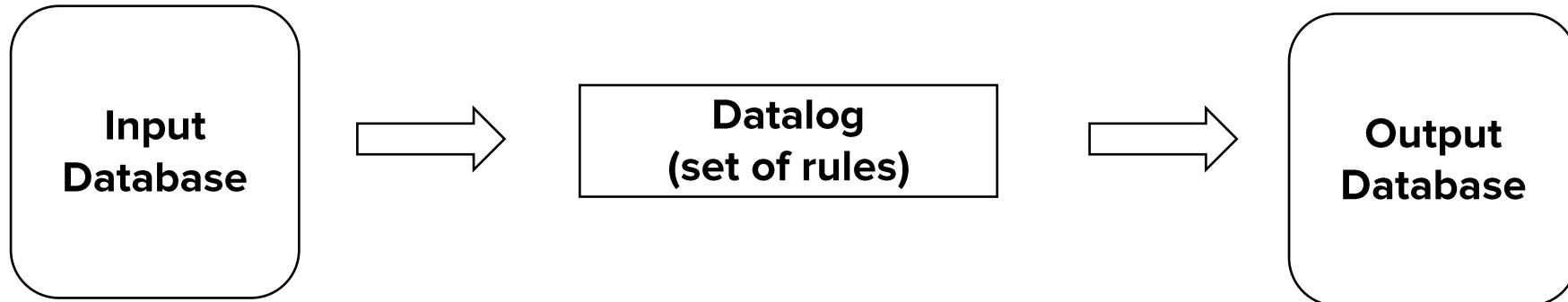
Results

Future Research Direction

# Datalog: Bottom-Up Logic Programming Language

| 4

A lightweight logic-programming language for deductive-database systems



Running the Datalog program extends data from input database creating the output database with all data transitively derivable via the program rules

- Ceri, S., Gottlob, G., & Tanca, L. (1989). What you always wanted to know about Datalog(and never dared to ask). *IEEE transactions on knowledge and data engineering*, 1(1), 146-166.
- Gilray, T., Kumar, S., & Micinski, K. (2021, March). Compiling data-parallel datalog. In *Proceedings of the 30th ACM SIGPLAN International Conference on Compiler Construction* (pp. 23-35).

# Classic Problems for Datalog

Transitive closure

Triangle counting

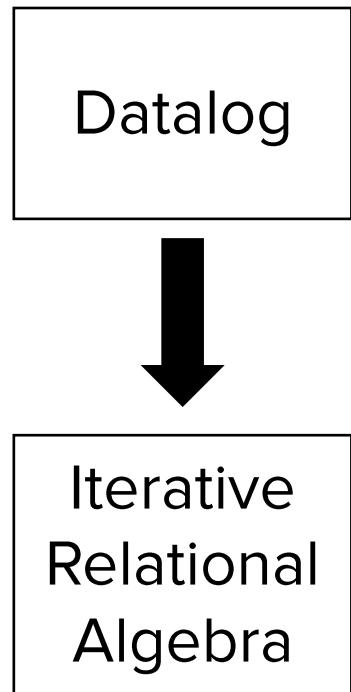
Finding maximal cliques

Finding frequent itemsets

Data mining

- Oege De Moor, Georg Gottlob, Tim Furche, and Andrew Sellers. Datalog Reloaded: First International Workshop, Datalog 2010, Oxford, UK, March 16–19, 2010. Revised Selected Papers, volume 6702. Springer, 2012.
- Jiwon Seo, Stephen Guo, and Monica S Lam. Socialite: Datalog extensions for efficient social network analysis. In 2013 IEEE 29th International Conference on Data Engineering (ICDE), pages 278–289. IEEE, 2013.

# Bottom-Up Logic Programming with Datalog



Datalog rule for computing **Transitive Closure (TC)**

$$\begin{aligned} T(x, y) &\leftarrow G(x, y) . \\ T(x, z) &\leftarrow T(x, y), G(y, z) . \end{aligned}$$



*Operationalized as a **fixed-point iteration** using  $F_G$*

$$F_G(T) \triangleq G \cup \Pi_{1,2}(\rho_{0/1}(T) \bowtie_1 G)$$

Relational algebra:



- Gilray, T., & Kumar, S. (2019, December). Distributed relational algebra at scale. In 2019 IEEE 26th International Conference on High Performance Computing, Data, and Analytics (HiPC) (pp. 12-22). IEEE.
- Kumar, S., & Gilray, T. (2020, June). Load-balancing parallel relational algebra. In International Conference on High Performance Computing (pp. 288-308). Springer, Cham.

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# Relational Algebra Primitives

- Main relational algebra primitives of two flat relations **R** and **S** are:
  - Union:  $R \cup S$
  - Intersection:  $R \cap S$
  - Cartesian product:  $R \times S$
  - Join:  $R \bowtie S$
  - Rename:  $\rho_{(i,j)}(R)$
  - Selection:  $\sigma_i(R)$
  - Projection:  $\Pi_{(i,j)}(R)$
- Differ from traditional set theory: **R** and **S** have a fixed arity

Relation		
UserID	UserName	UserEmail
101	Alice	alice@example.com
102	Bob	bob@example.com

Attribute  
 (Column)

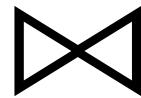
Tuple  
 (Row)

- Sidharth Kumar and Thomas Gilray. Distributed relational algebra at scale. In International Conference on High Performance Computing, Data, and Analytics (HiPC). IEEE, 2019.
- Sidharth Kumar and Thomas Gilray. Load-balancing parallel relational algebra. In International Conference on High Performance Computing, pages 288–308. Springer, 2020.

# Example of Natural Join

User

<b>UserID</b>	<b>UserName</b>	<b>UserEmail</b>
101	Alice	alice@example.com
102	Bob	bob@example.com
103	Eve	eve@example.com



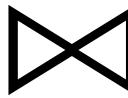
Order

<b>UserID</b>	<b>OrderTotal</b>	<b>Items</b>
101	25.69	2
102	145.66	3
103	12.11	1
103	44.00	2

# Example of Natural Join

User

UserID	UserName	UserEmail
101	Alice	alice@example.com
102	Bob	bob@example.com
103	Eve	eve@example.com



Order

UserID	OrderTotal	Items
101	25.69	2
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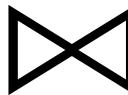


UserID	UserName	UserEmail	OrderTotal	Items
101	Alice	alice@example.com	25.69	2

# Example of Natural Join

User

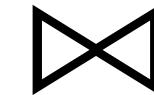
UserID	UserName	UserEmail
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Order

UserID	OrderTotal	Items
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User



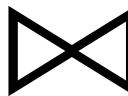
Order

UserID	UserName	UserEmail	OrderTotal	Items
101	Alice	alice@example.com	25.69	2
102	Bob	bob@example.com	145.66	3

# Example of Natural Join

User

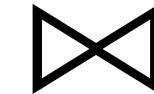
<b>UserID</b>	<b>UserName</b>	<b>UserEmail</b>
101	Alice	alice@example.com
102	Bob	bob@example.com
103	Eve	eve@example.com



Order

<b>UserID</b>	<b>OrderTotal</b>	<b>Items</b>
101	25.69	2
102	145.66	3
103	12.11	1
103	44.00	2

User



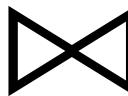
Order

<b>UserID</b>	<b>UserName</b>	<b>UserEmail</b>	<b>OrderTotal</b>	<b>Items</b>
101	Alice	alice@example.com	25.69	2
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103	Eve	eve@example.com	12.11	1

# Example of Natural Join

User

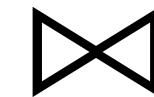
<b>UserID</b>	<b>UserName</b>	<b>UserEmail</b>
101	Alice	alice@example.com
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103	Eve	eve@example.com



Order

<b>UserID</b>	<b>OrderTotal</b>	<b>Items</b>
101	25.69	2
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User



Order

<b>UserID</b>	<b>UserName</b>	<b>UserEmail</b>	<b>OrderTotal</b>	<b>Items</b>
101	Alice	alice@example.com	25.69	2
102	Bob	bob@example.com	145.66	3
103	Eve	eve@example.com	12.11	1
103	Eve	eve@example.com	44.00	2

# Duplicates on Join Result

User  $\bowtie$  Order

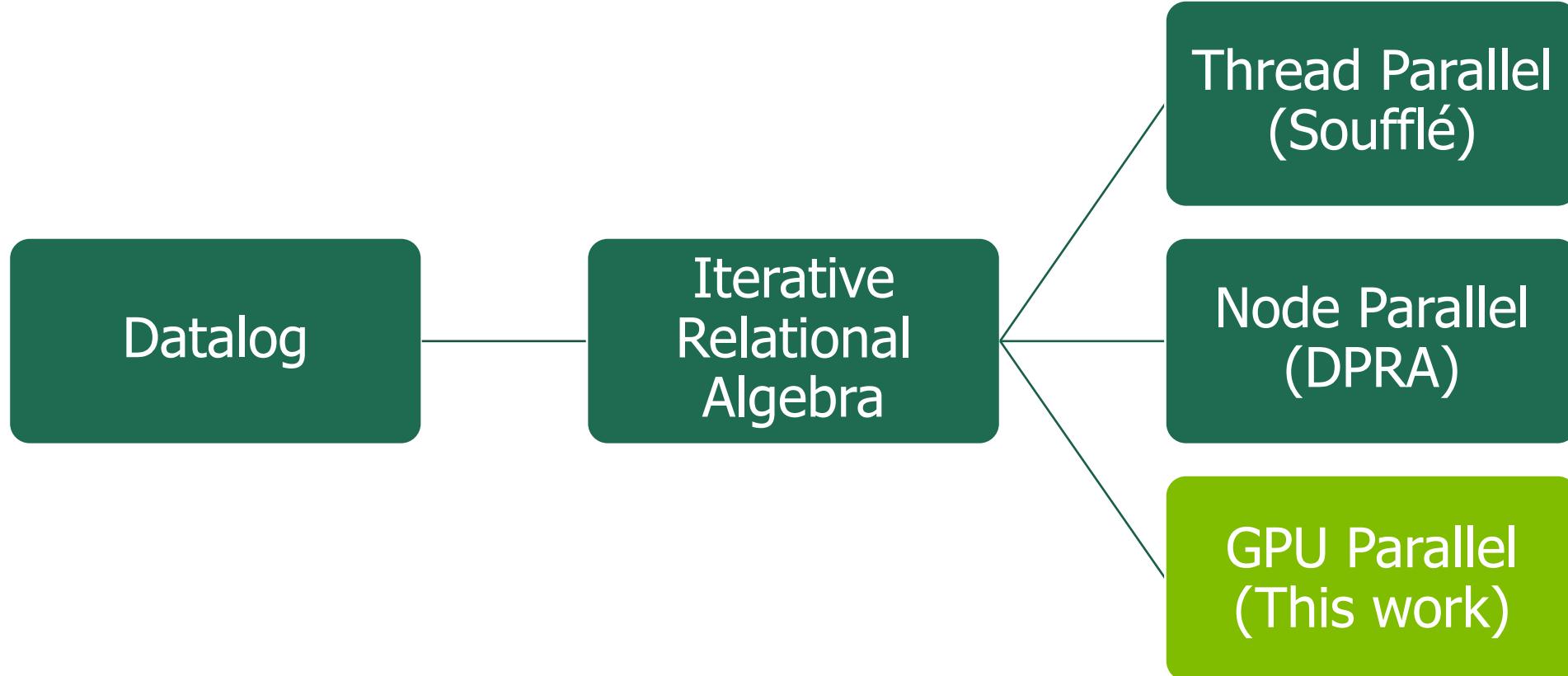
UserID	UserName	UserEmail	OrderTotal	Items
101	Alice	alice@example.com	25.69	2
102	Bob	bob@example.com	145.66	3
103	Eve	eve@example.com	12.11	1
103	Eve	eve@example.com	44.00	2

$\Pi_{(\text{UserName}, \text{UserEmail})}(\text{User} \bowtie \text{Order})$

UserName	UserEmail
Alice	alice@example.com
Bob	bob@example.com
Eve	eve@example.com
Eve	eve@example.com

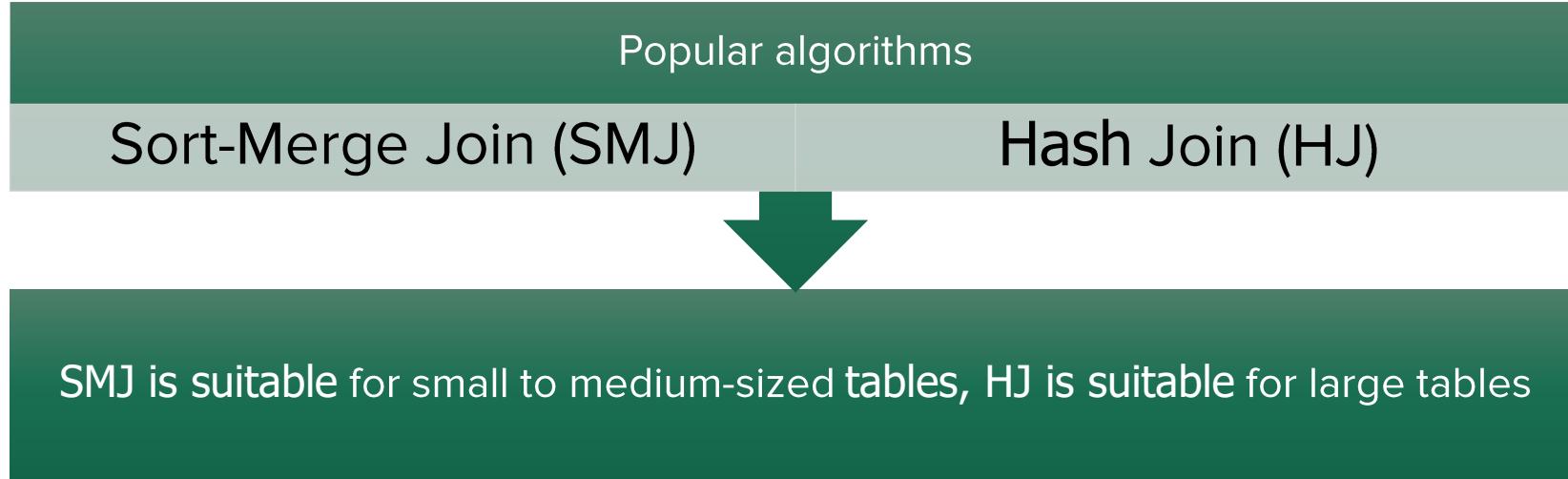
# Towards Parallel Relational Algebra

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- Herbert Jordan, Bernhard Scholz, and Pavle Subotić. Soufflé: On synthesis of program analyzers. In International Conference on Computer Aided Verification, pages 422–430. Springer, 2016.
- Kumar, S., & Gilray, T. (2019). Distributed relational algebra at scale. In International Conference on High Performance Computing, Data, and Analytics (HiPC). IEEE (Vol. 1).
- Thomas Gilray, Sidharth Kumar, and Kristopher Micinski. Compiling data-parallel datalog. In Proceedings of the 30th ACM SIGPLAN International Conference on Compiler Construction, CC 2021, page 23–35, New York, NY, USA, 2021. Association for Computing Machinery.

# Parallel Join: Algorithms



- Chengxin Guo, Hong Chen, Feng Zhang, and Cuiping Li. Parallel hybrid join algorithm on gpu. 2019 IEEE 21st International Conference on High Performance Computing and Communications; IEEE17th International Conference on Smart City; IEEE 5th International Conference on Data Science andSystems (HPCC/SmartCity/DSS), pages 1572–1579, 2019.
- Hongzhi Wang, Ning Li, Zheng ke Wang, and Jianing Li. Gpu-based efficient join algorithms on hadoop. The Journal of Supercomputing, 77:292 – 321, 2020.

# Research Gaps in Parallel Join Implementations



GPU-based join implementations does not sort result (by default)



Challenge for iterated relational algebra algorithms



Negative impact on algorithm performance



Memory overhead in Python libraries

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**Transitive Closure Computation**

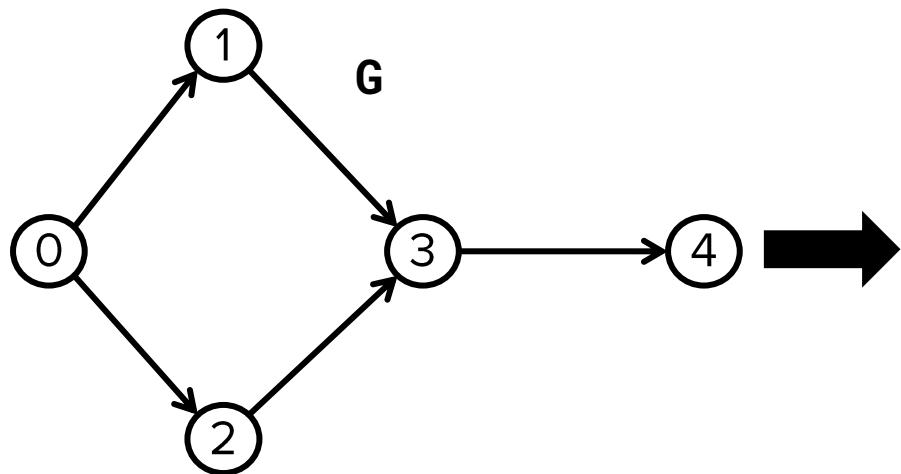
Experimental Setup & Dataset

Results

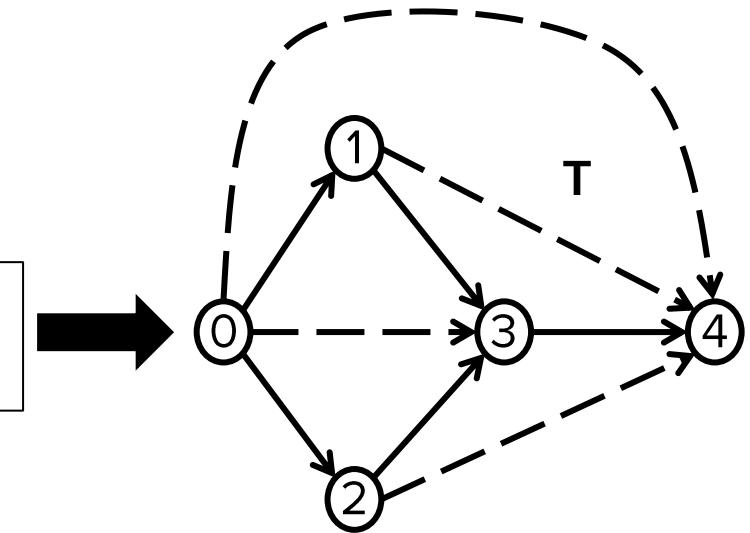
Future Research Direction

# Transitive Closure: Logical Inference for Graphs

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$T(x, y) \leftarrow G(x, y)$   
 $T(x, z) \leftarrow T(x, y), G(y, z)$



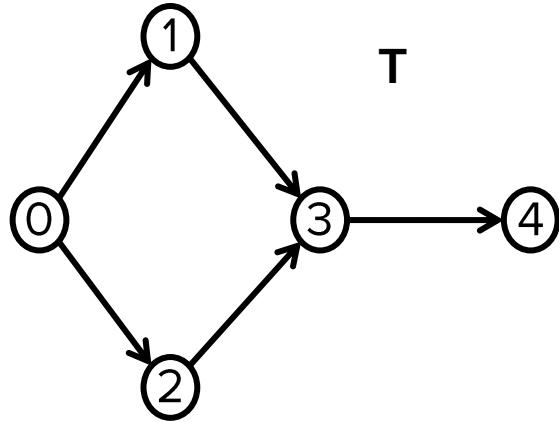
0	1
0	2
1	3
2	3
3	4

0	1
0	2
1	3
2	3
3	4
0	3
1	4
2	4
0	4

$$F_G(T) \triangleq G \cup \Pi_{1,2}(\rho_{0/1}(T) \bowtie_1 G)$$

| ↗

# Transitive Closure: Iterations

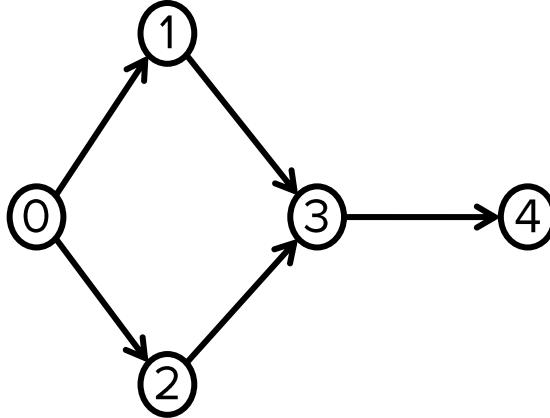


0	1
0	2
1	3
2	3
3	4

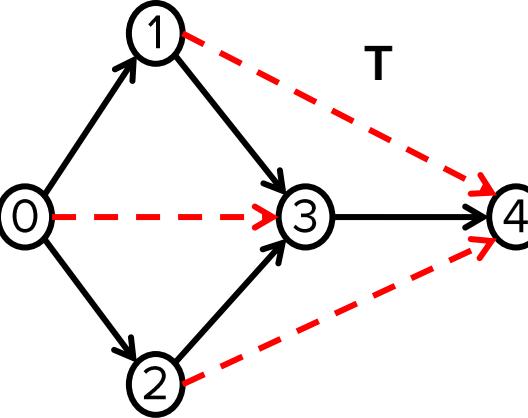
$$F_G(T) \triangleq G \cup \Pi_{1,2}(\rho_{0/1}(T) \bowtie_1 G)$$

# Transitive Closure: Iterations 1

| ↗



0	1
0	2
1	3
2	3
3	4

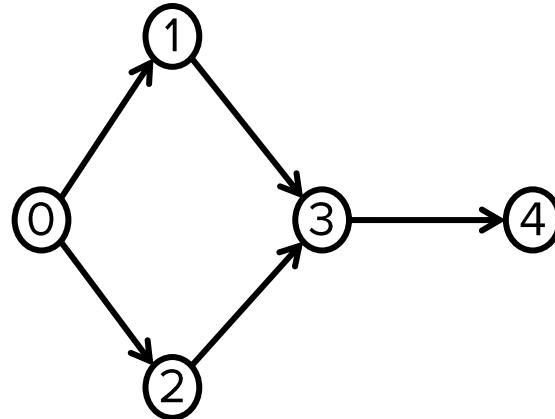


0	1
0	2
1	3
2	3
3	4
0	3
1	4
2	4

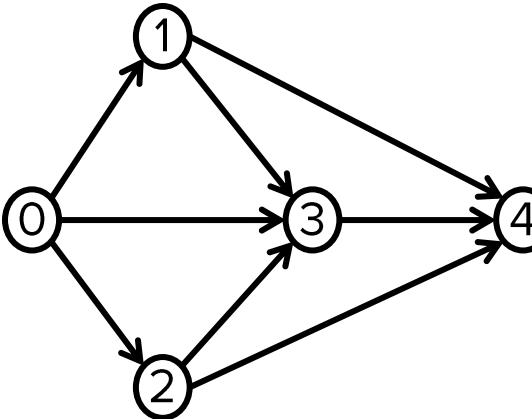
$$F_G(T) \triangleq G \cup \Pi_{1,2}(\rho_{0/1}(T) \bowtie_1 G)$$

# Transitive Closure: Iterations 2

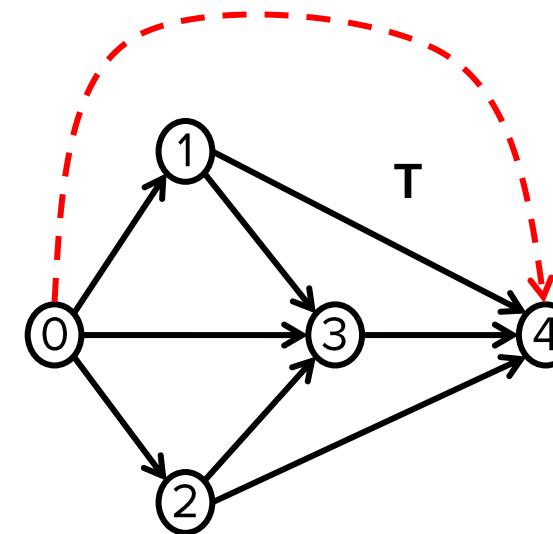
| ↴



0	1
0	2
1	3
2	3
3	4



0	1
0	2
1	3
2	3
3	4
0	3
1	4
2	4

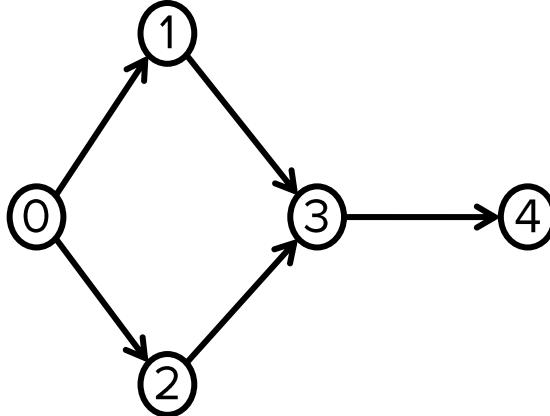


0	1
0	2
1	3
2	3
3	4
0	3
1	4
2	4
0	4

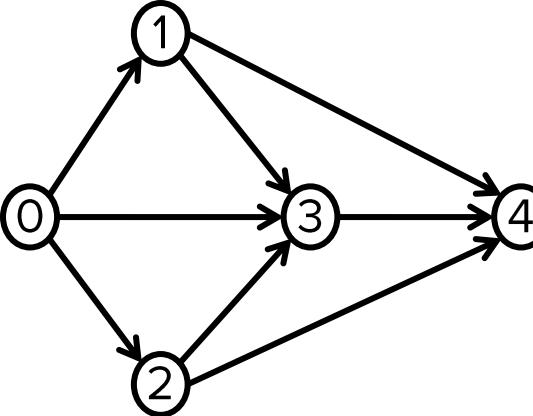
$$F_G(T) \triangleq G \cup \Pi_{1,2}(\rho_{0/1}(T) \bowtie_1 G)$$

# Transitive Closure: Iterations 3

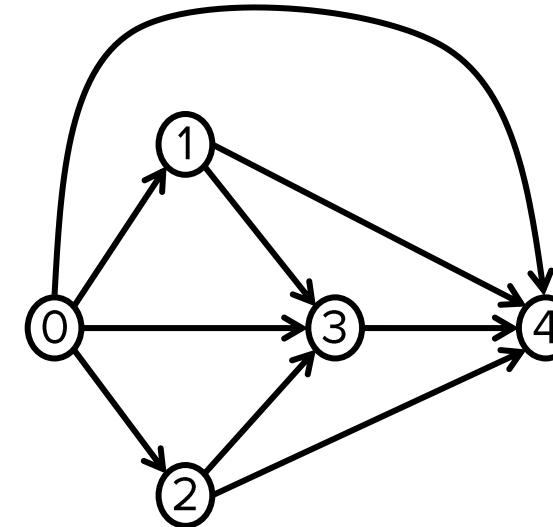
| 23



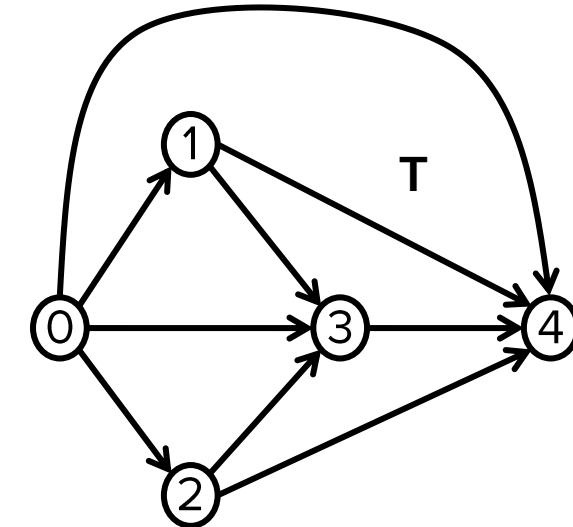
0	1
0	2
1	3
2	3
3	4



0	1
0	2
1	3
2	3
3	4
0	3
1	4
2	4



0	1
0	2
1	3
2	3
3	4
0	3
1	4
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1	3
2	3
3	4
0	3
1	4
2	4
0	4

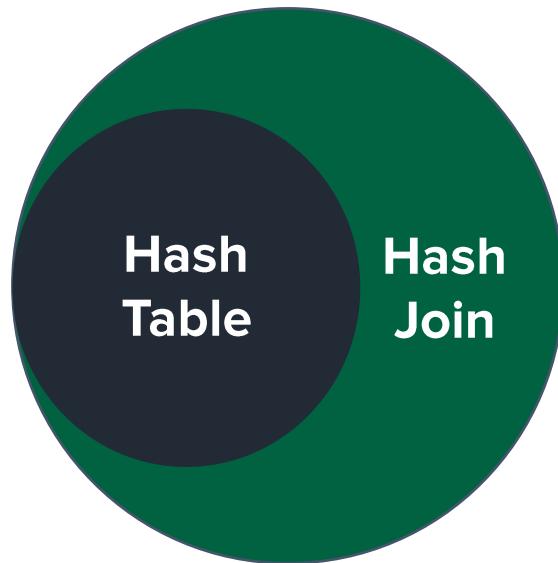
# TC Computation in Iterated Relational Algebra

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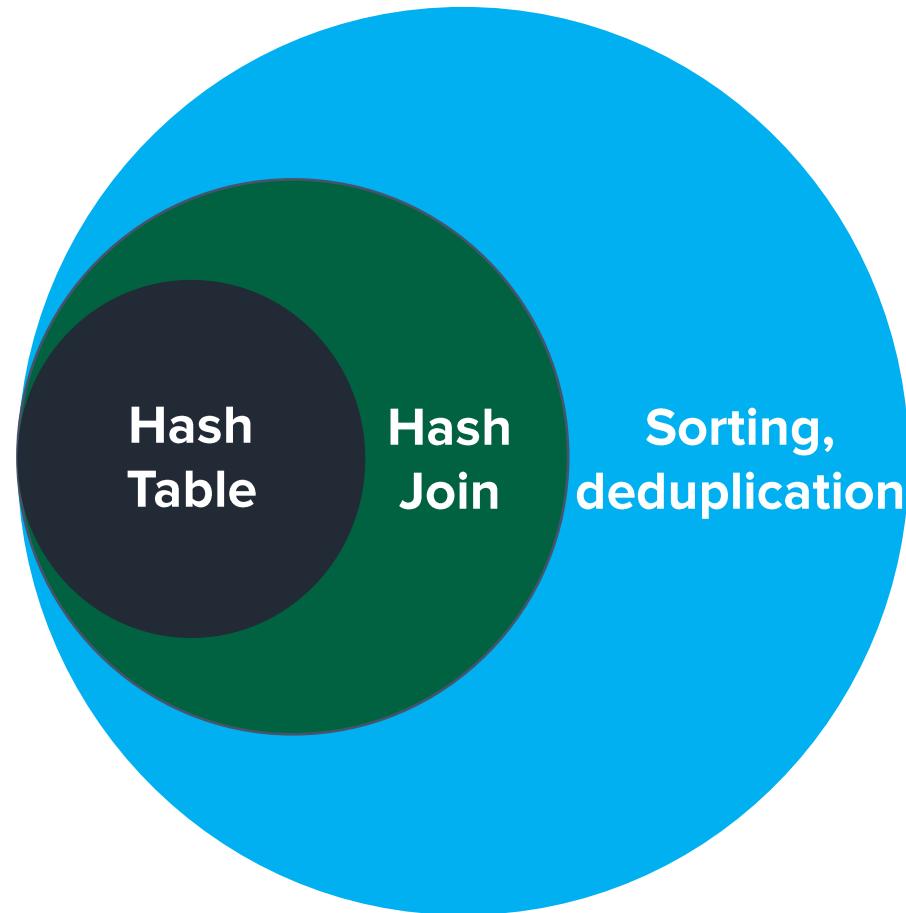
# TC Computation in Iterated Relational Algebra

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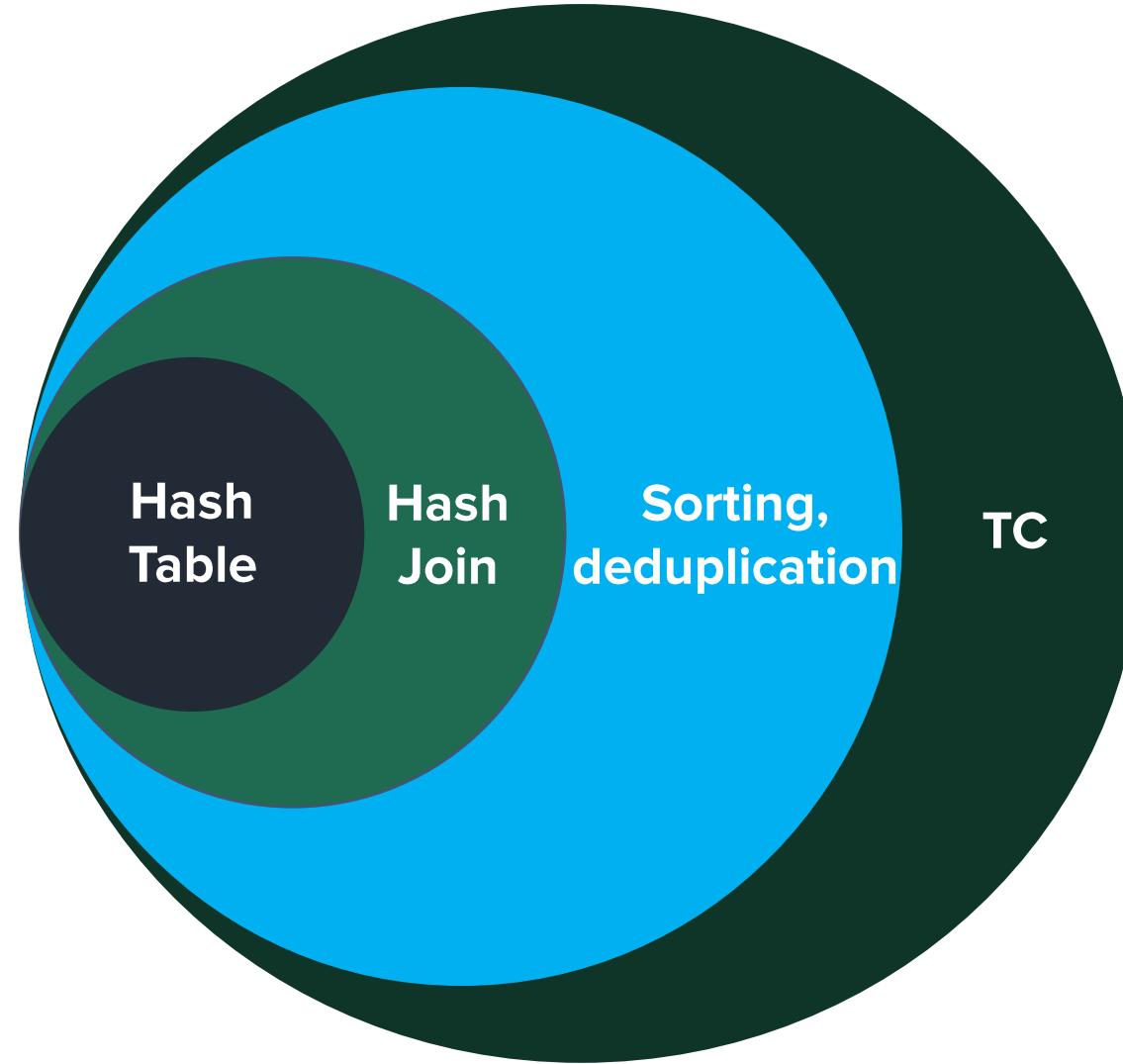
# TC Computation in Iterated Relational Algebra

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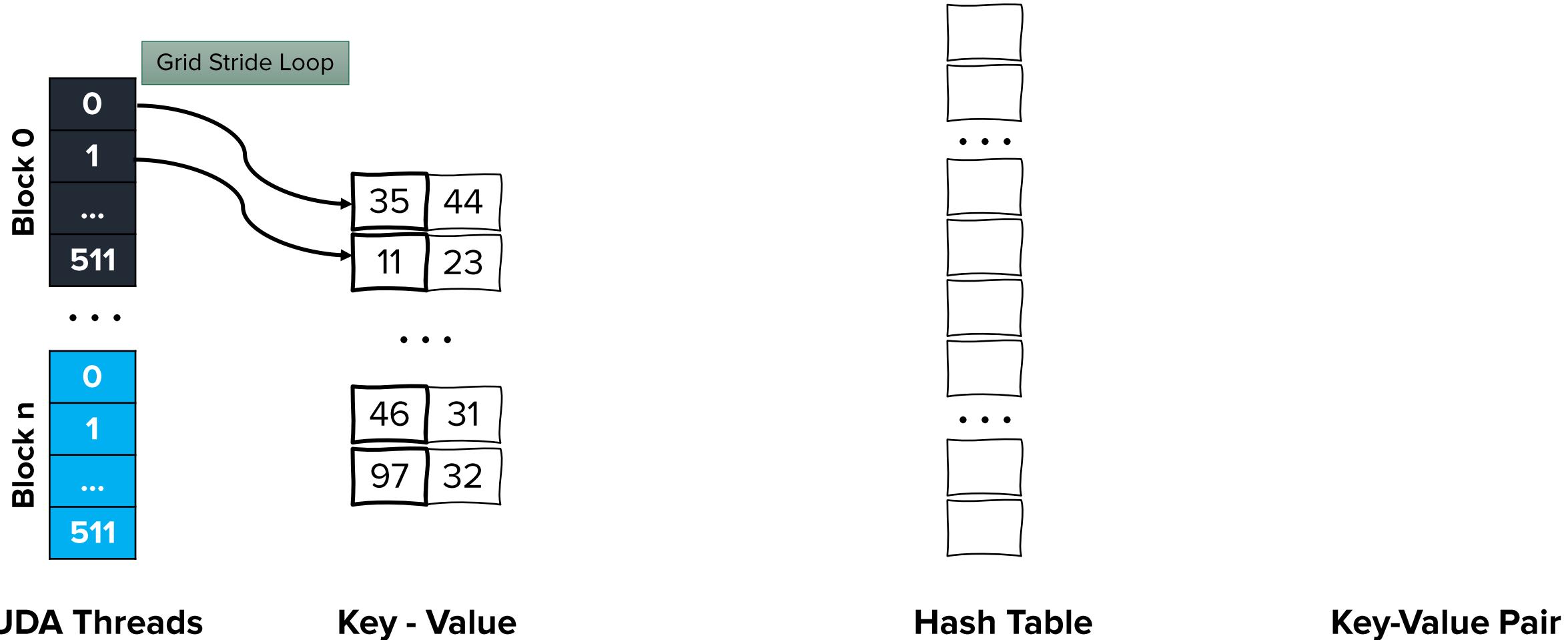
# TC Computation in Iterated Relational Algebra

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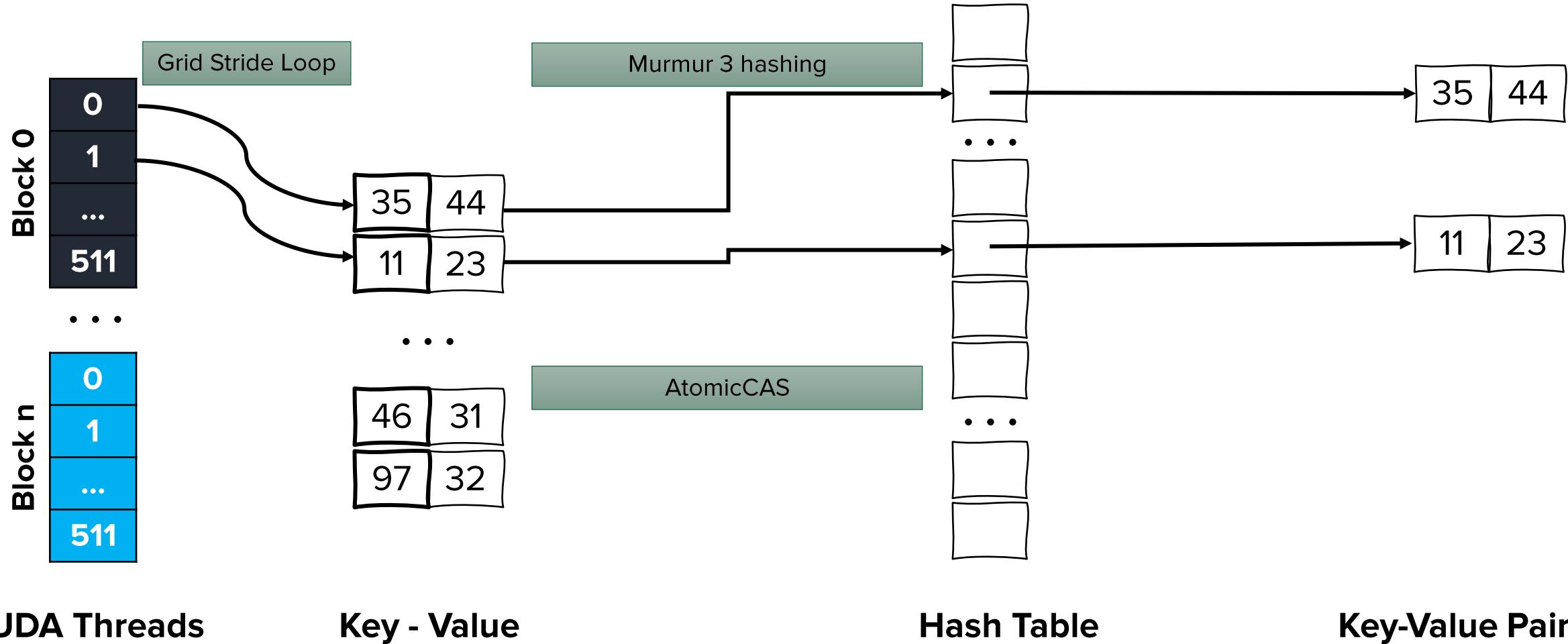


# Hash Table (Open Addressing, Linear Probing)

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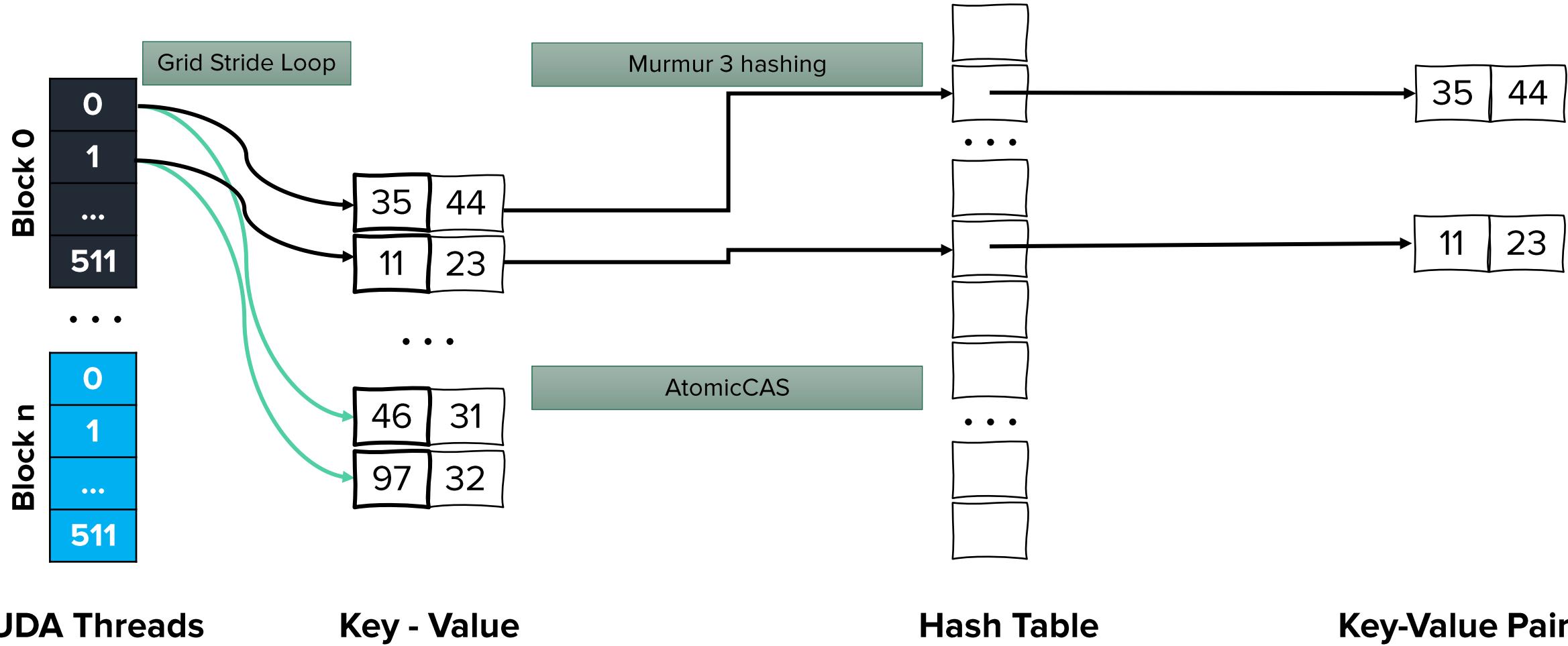


# Hash Table (Open Addressing, Linear Probing)



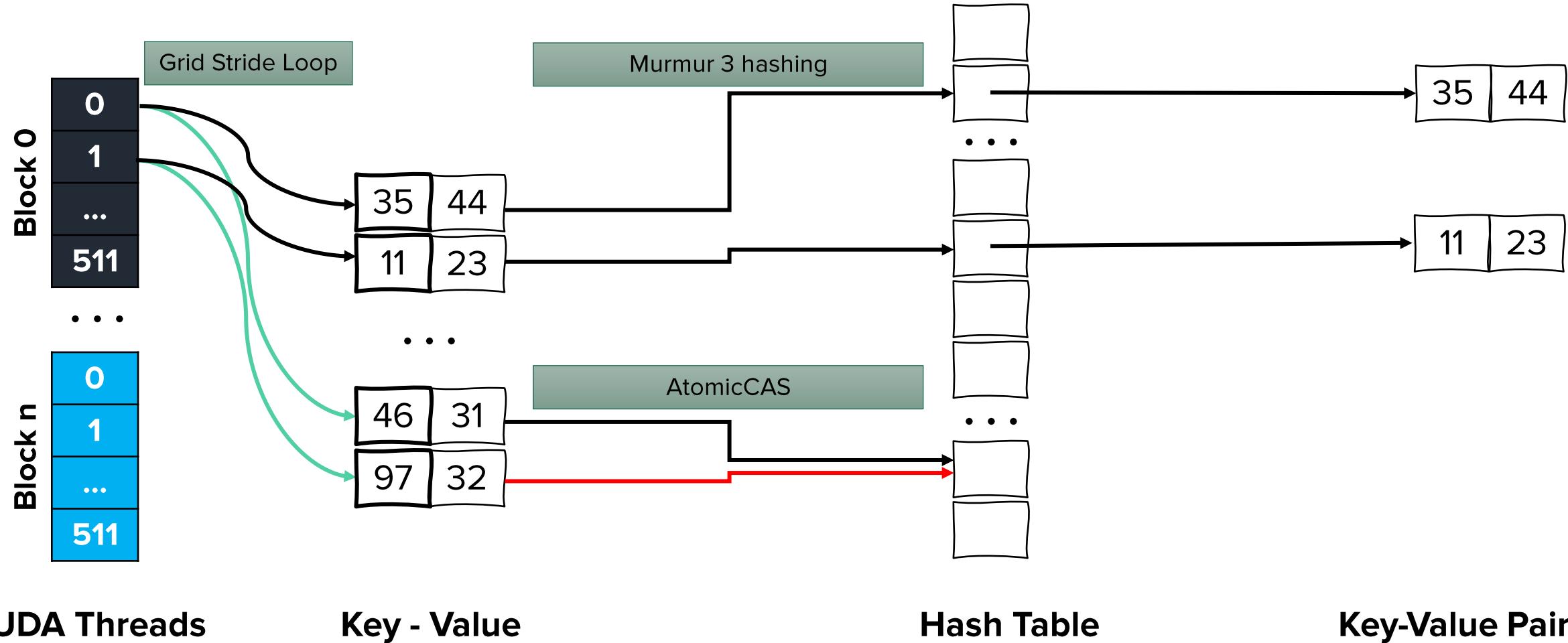
# Hash Table (Open Addressing, Linear Probing)

| 30



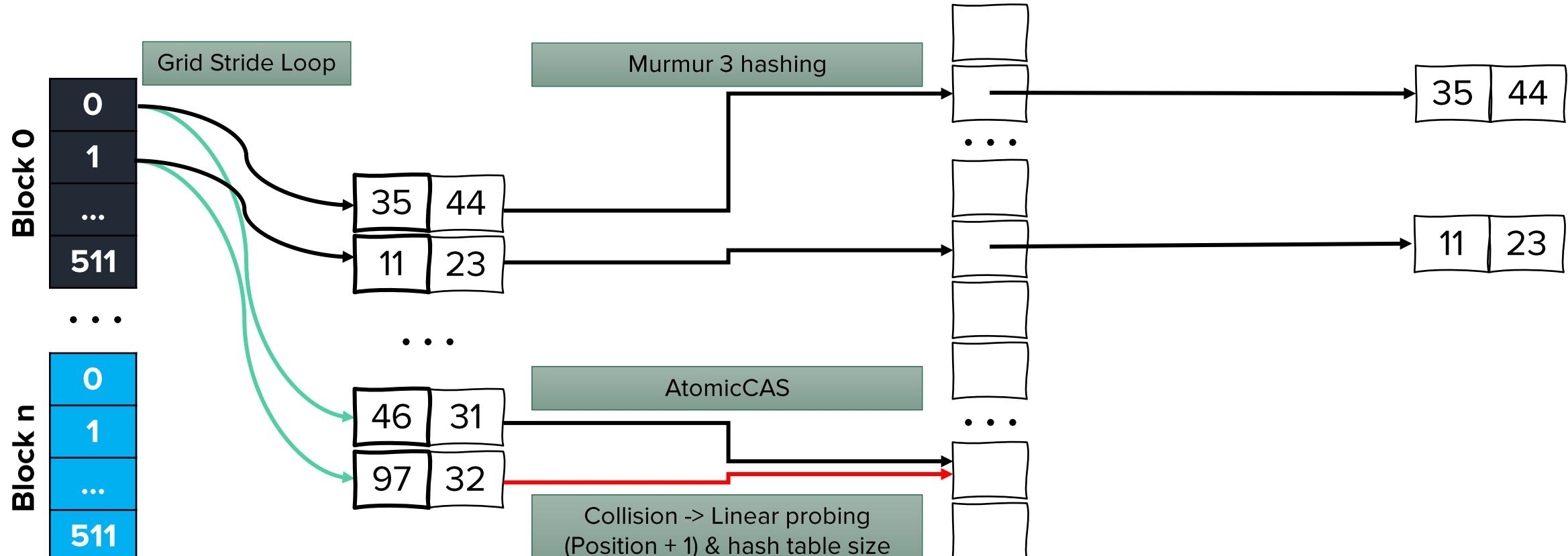
# Hash Table (Open Addressing, Linear Probing)

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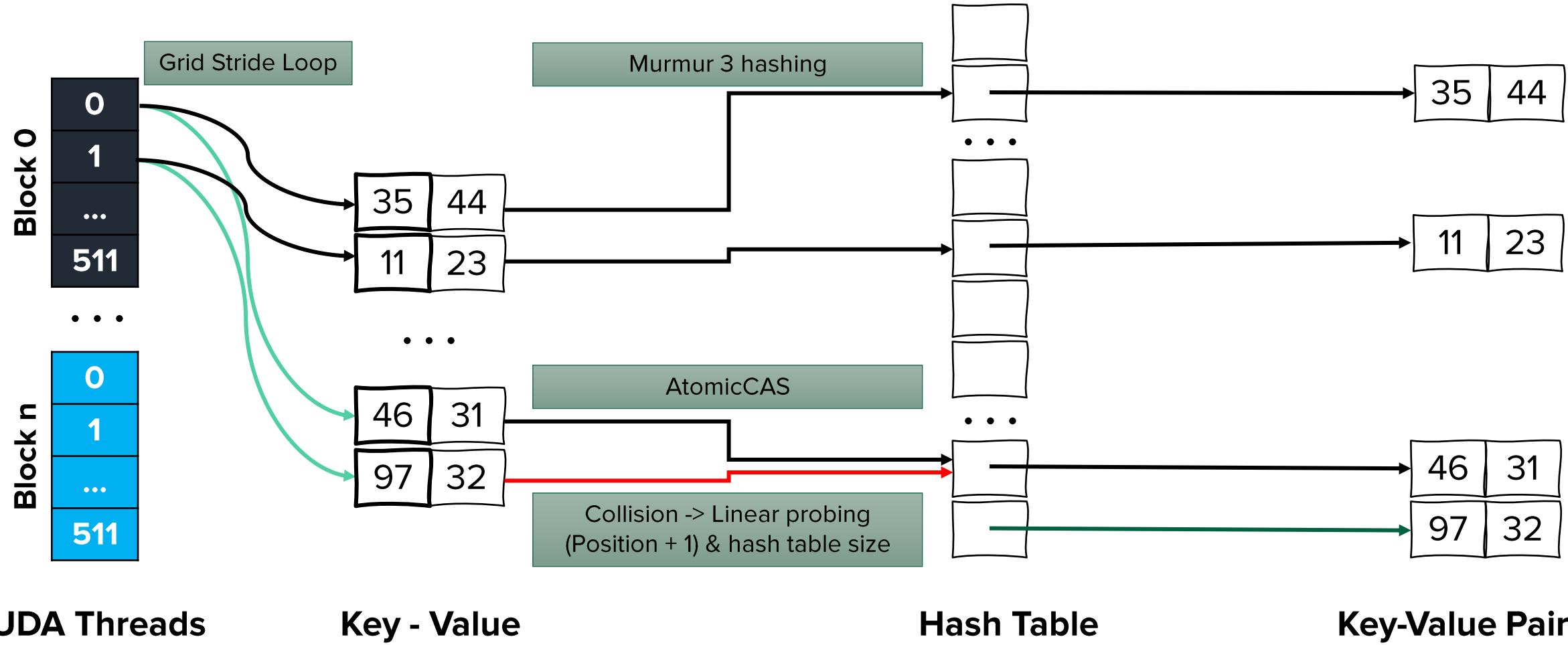
# Hash Table (Open Addressing, Linear Probing)

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# Hash Table (Open Addressing, Linear Probing)

33



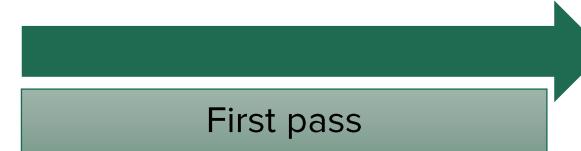
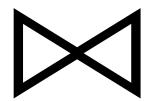
# Performing Hash Join on GPU

Static Hash Table

Key	Value

Reverse Relation

Key	Value



Calculate join size

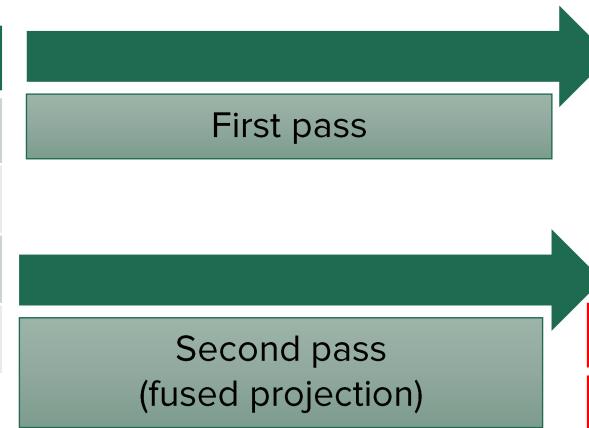
# Performing Hash Join on GPU

Static Hash Table

Key	Value

Reverse Relation

Key	Value



Calculate join size

Join Result

Join Key	Value 1	Value 2
X		

Prefix sum

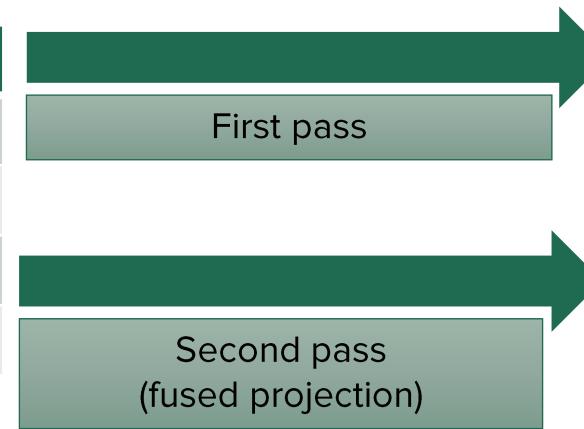
# Performing Hash Join on GPU

Static Hash Table

Key	Value

Reverse Relation

Key	Value



Calculate join size

Prefix sum

Join Result

Join Key	Value 1	Value 2
X		
X		
X		
X		
X		
X		
X		
X		

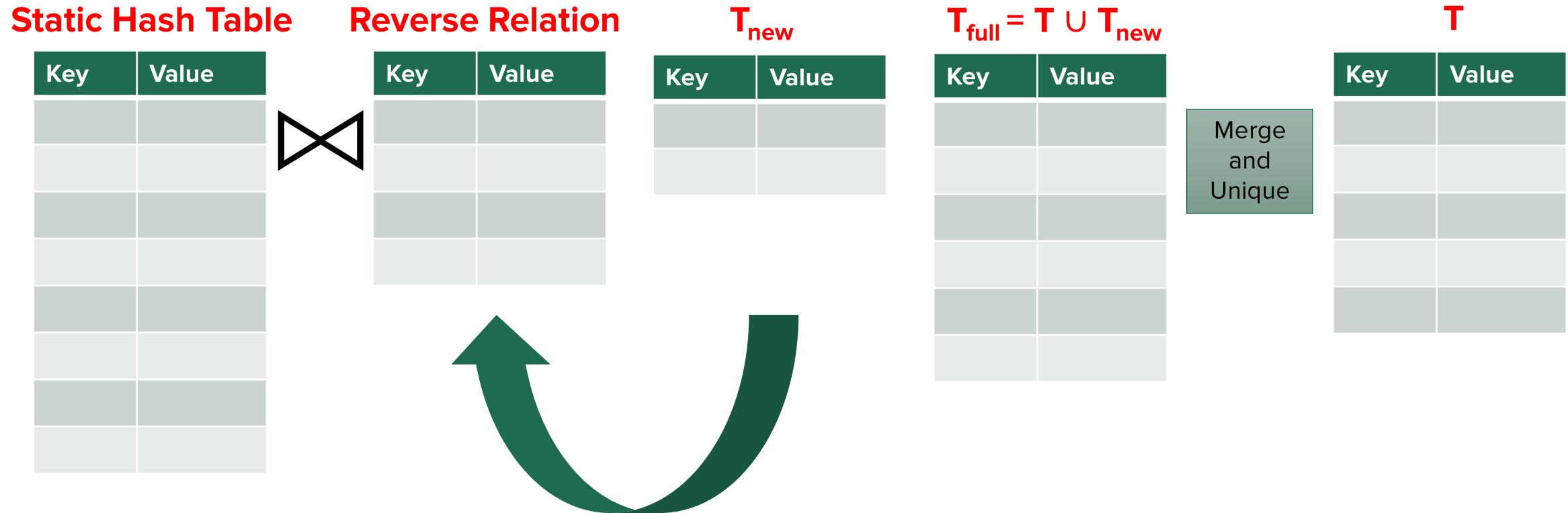
Sort and Unique

Deduplicated Join Result

Key	Value

## Transitive closure computation (single iteration)

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Process continues until there is no new facts are discovered in an iteration

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# Experiment Platform and Datasets

ThetaGPU supercomputer from Argonne National Lab

CPU: AMD EPYC 7742 processors with 3.31GHz clock speed, 128 cores

## GPU

- NVIDIA A100 Tensor Core GPU with 40GB GPU memory
- 108 multiprocessors on device (SM)

## Environment

- CUDA version 11.4, 3,456 x 512 (blocks per grid x threads per block)
- Souffle version 2.3 with 128 threads
- cuDF package inside conda environment

## Datasets

- Stanford large network dataset collection
  - SuiteSparse matrix collection
  - Road network real datasets collection
- 
- Leskovec, J., & Krevl, A. (2014). SNAP Datasets: Stanford large network dataset collection.

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**Results**

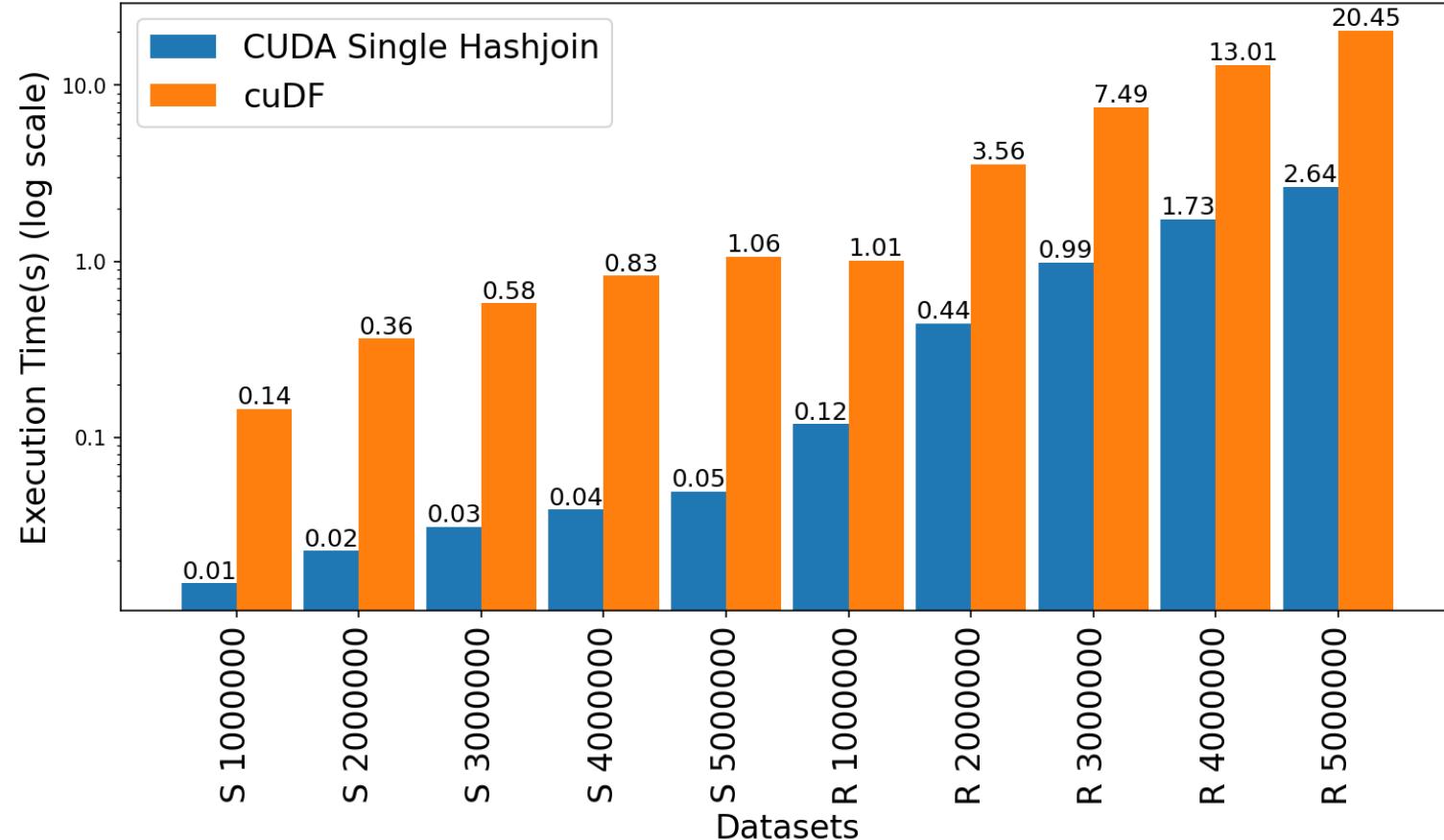
Future Research Direction

# Hash Table Performance

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- Build rate:
  - Random synthetic graph: 400 million keys/second
  - String graph: 4 billion keys/second
- Load factors are varied to ensure less memory overhead

# Join Performance Comparison: CUDA vs cuDF



- Leadership Computing Facility, A. (2022). Argonne Leadership Computing Facility. Theta GPU Nodes. URL: <https://www.alcf.anl.gov/support-center/theta-gpu-nodes>

# CUDA Advantages over Dataframe

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Fuse operations

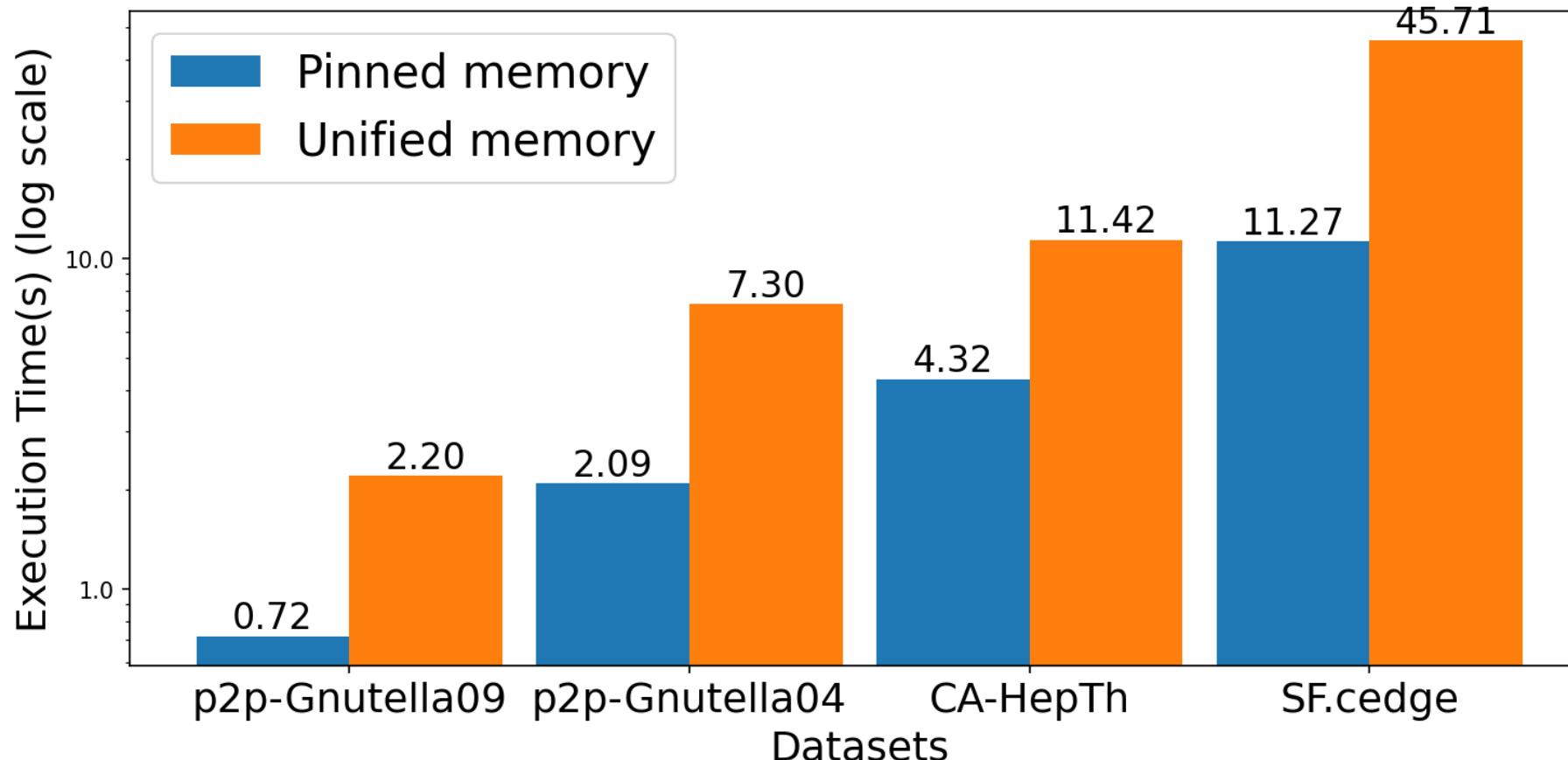
Thread-block configuration

Memory management

Optimize data structure

- Jason. Sanders. CUDA by example : an introduction to general-purpose GPU programming. AddisonWesley, Upper Saddle River, NJ, 2011.
- John Cheng, Max Grossman, and Ty McKercher. Professional CUDA c programming. John Wiley & Sons, 2014

# TC Performance Comparison: Memory Schemes



- Leadership Computing Facility, A. (2022). Argonne Leadership Computing Facility. Theta GPU Nodes. URL: <https://www.alcf.anl.gov/support-center/theta-gpu-nodes>

# TC Performance Comparison: CUDA vs Soufflé vs cuDF

Dataset	Type	Rows	TC size	Iterations	CUDA Hashjoin(s)	Soufflé(s)	cuDF(s)
fe_ocean	U	409,593	1,669,750,513	247	138.237	536.233	Out of Memory
p2p-Gnutella31	D	147,892	884,179,859	31	Out of Memory	128.917	Out of memory
usroads	U	165,435	871,365,688	606	364.554	222.761	Out of Memory
fe_body	U	163,734	156,120,489	188	47.758	29.07	Out of Memory
loc-Brightkite	U	214,078	138,269,412	24	15.88	29.184	Out of Memory
SF.cedge	U	223,001	80,498,014	287	11.274	17.073	64.417
fe_sphere	U	49,152	78,557,912	188	13.159	20.008	80.077
CA-HepTh	D	51,971	74,619,885	18	4.318	15.206	26.115
p2p-Gnutella04	D	39,994	47,059,527	26	2.092	7.537	14.005
p2p-Gnutella09	D	26,013	21,402,960	20	0.72	3.094	3.906
wiki-Vote	D	103,689	11,947,132	10	1.137	3.172	6.841
cti	U	48,232	6,859,653	53	0.295	1.496	3.181
delaunay_n16	U	196,575	6,137,959	101	1.137	1.612	5.596
luxembourg_osm	U	119,666	5,022,084	426	1.322	2.548	8.194
ego-Facebook	U	88,234	2,508,102	17	0.544	0.606	3.719
cal.cedge	U	21,693	501,755	195	0.489	0.455	2.756
TG.cedge	U	23,874	481,121	58	0.198	0.219	0.857
wing	U	121,544	329,438	11	0.085	0.193	0.905
OL.cedge	U	7,035	146,120	64	0.148	0.181	0.523

# Cases Where Souffle Outperforms CUDA

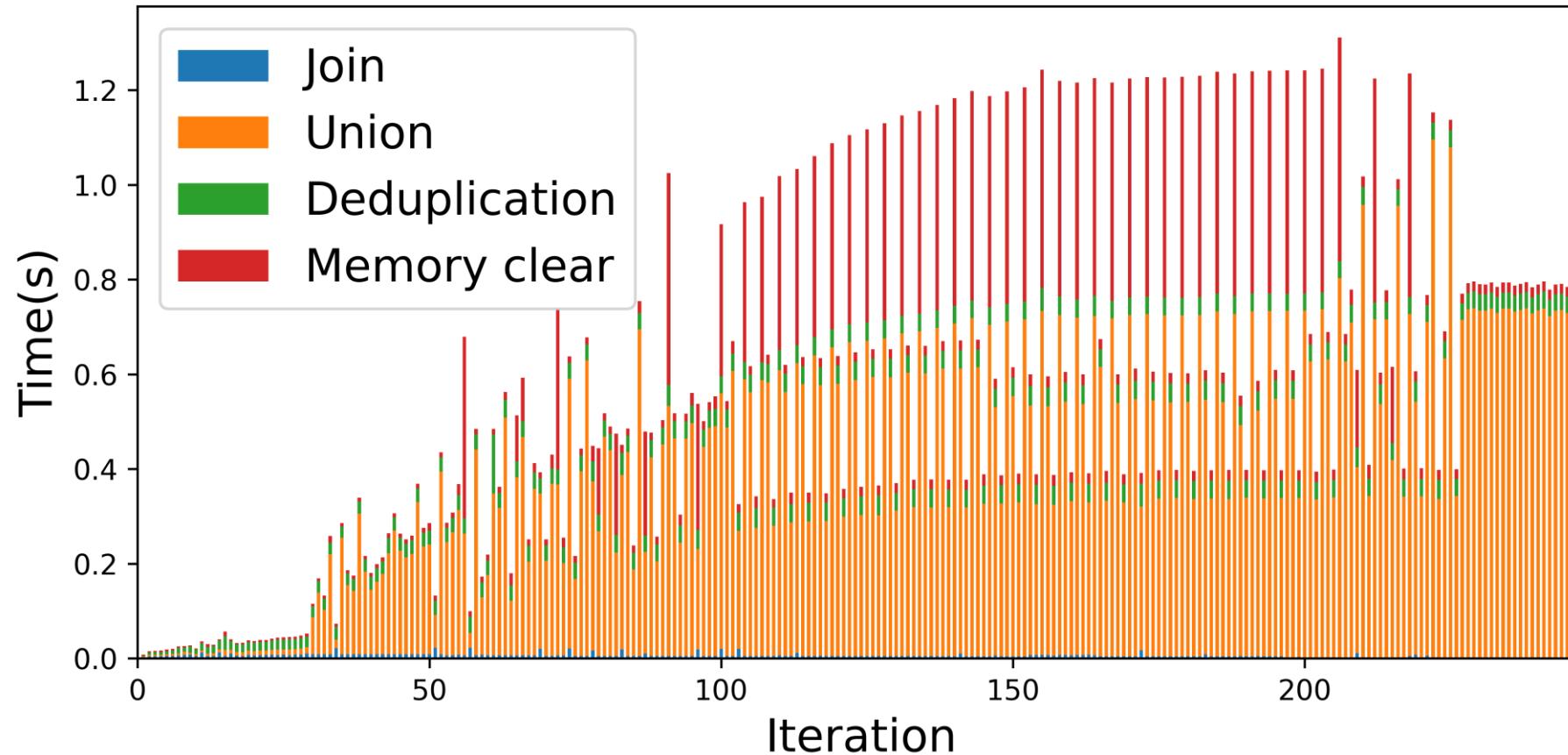
| 46



Overflows GPU memory when  
higher workload/iteration

Underperforms when less  
work for GPU/iteration

# Operations Breakdown per Iteration (fe\_ocean)



- Leadership Computing Facility, A. (2022). Argonne Leadership Computing Facility. Theta GPU Nodes. URL: <https://www.alcf.anl.gov/support-center/theta-gpu-nodes>

# Contributions

High Performance GPU hash table for iterative RA

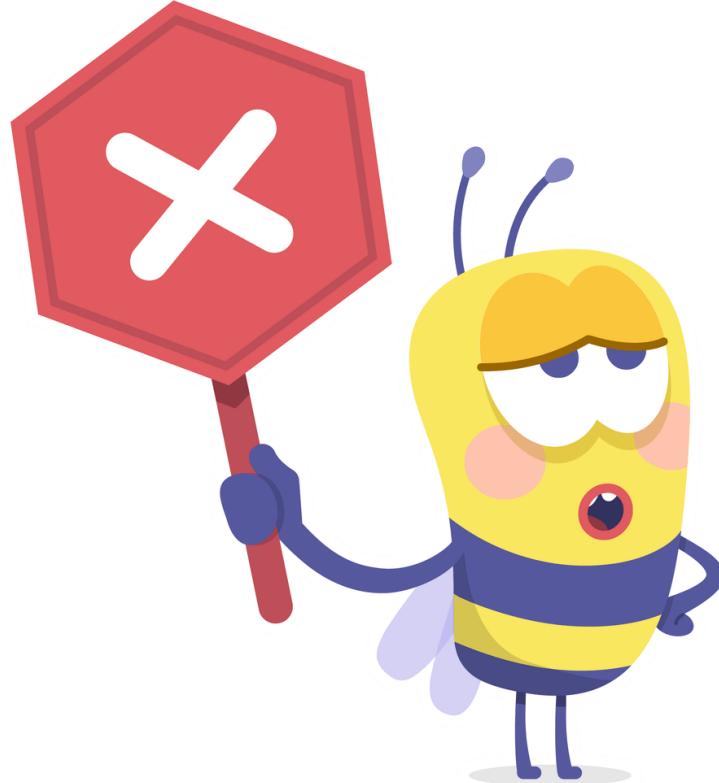
Operations optimization (fuse join and projection)

Overcome deduplication challenge

Efficient GPU memory management (pinned and buffer clearance)

# Limitations

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---

Limited to a single GPU that dictates scaling by available VRAM on the GPU

---

Memory overflow error for larger graphs

---

Open addressing based hash table causes memory overhead

---

# Table of Contents

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Datalog

Iterative Relational Algebra on GPU

Transitive Closure Computation

Experimental Setup & Dataset

Results

**Future Research Direction**

# Future Work

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Develop

Multi-node multi-GPU backend for Datalog to perform iterated relational algebra operations tailored for GPU

Compare

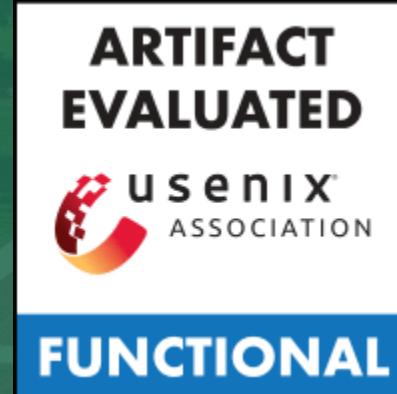
Different Parallel Programming Models performance on iterative relational join

Extend

State-of-the-art multi-node CPU-based Datalog-like language SLOG to leverage our GPU-based solutions



<https://github.com/harp-lab/usenixatc23>



# Thank you!

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## HARP Lab

High-performance Automated Reasoning and Programming Lab

<https://github.com/harp-lab/>



# Appendix

# DataFrame Based Datalog Applications

## ✓ Advantages

- ✓ Abstract memory management
- ✓ Abstract thread block configuration
- ✓ Same API signatures for CPU and GPU
- ✓ Easy-to-code interface

## ✗ Limitations

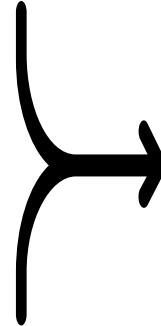
- ✗ No fusing
- ✗ Memory and computation overhead
- ✗ No consecutive operation
- ✗ Memory limitation

- A. R. Shovon, L. R. Dyken, O. Green, T. Gilray and S. Kumar, "Accelerating Datalog applications with cuDF," 2022 IEEE/ACM Workshop on Irregular Applications: Architectures and Algorithms (IA3), Dallas, TX, USA, 2022, pp. 41-45
- Green, O., Du, Z., Patel, S., Xie, Z., Liu, H., & Bader, D. A. (2021, December). Anti-Section Transitive Closure. In 2021 IEEE 28th International Conference on High Performance Computing, Data, and Analytics (HiPC) (pp. 192-201). IEEE.
- Team, R. D. (2018). RAPIDS: Collection of libraries for end to end GPU data science. NVIDIA, Santa Clara, CA, USA. <https://rapids.ai>

# Datalog Example

**Facts:**

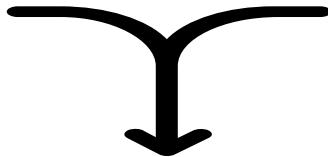
$\text{parents}(x,y).$   
 $\text{children}(y,x).$



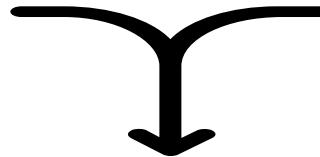
*extensional*

**Rules:**

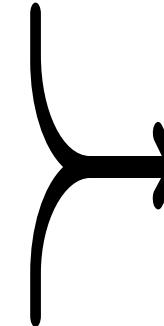
$\text{grandparent}(x,y) :- \text{parents}(x,z), \text{parents}(z,y).$



*head*



*body*



*intensional*

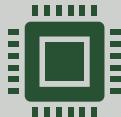
- Michael Stonebraker. Readings in database systems. Morgan Kaufmann Publishers Inc., 1988
- Evgeny Dantsin, Thomas Eiter, Georg Gottlob, and Andrei Voronkov. Complexity and expressive power of logic programming. *ACM Comput. Surv.*, 33(3):374–425, sep 2001.
- David Maier, K Tuncay Tekle, Michael Kifer, and David S Warren. Datalog: concepts, history, and outlook. In *Declarative Logic Programming: Theory, Systems, and Applications*, pages 3–100. 2018.

# Parallel Join

57



**What:** Perform relational join operation simultaneously on a number of processors or machines



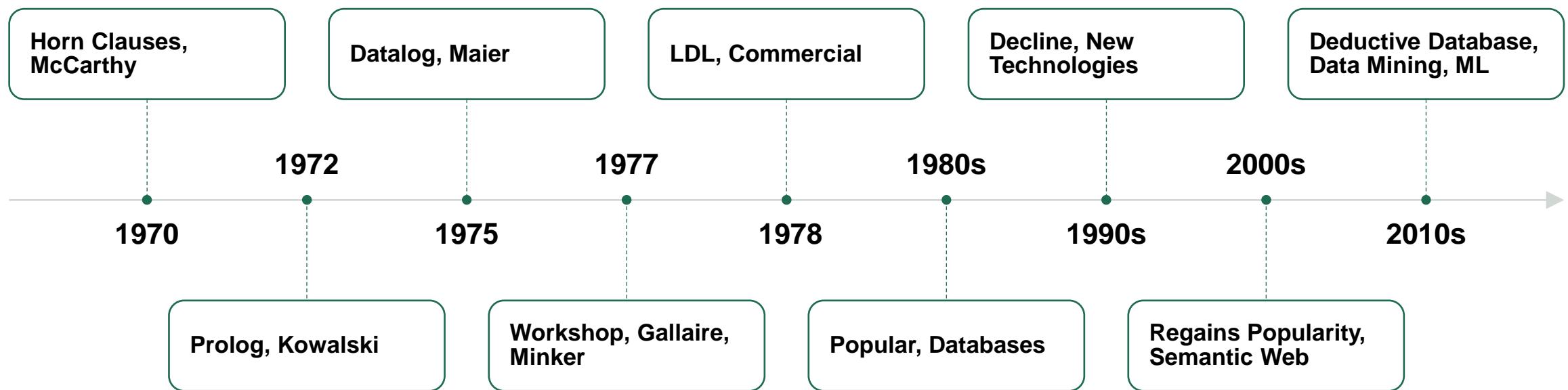
**When:** Useful when input data is enormous and the join is computationally costly



**How:** Divide the data into partitions and assign each partition to a different processor

- Daniel Zinn, Haicheng Wu, Jin Wang, Molham Aref, and Sudhakar Yalamanchili. General-purpose join algorithms for large graph triangle listing on heterogeneous systems. In Proceedings of the 9th Annual Workshop on General Purpose Processing Using Graphics Processing Unit, pages 12–21, 2016.

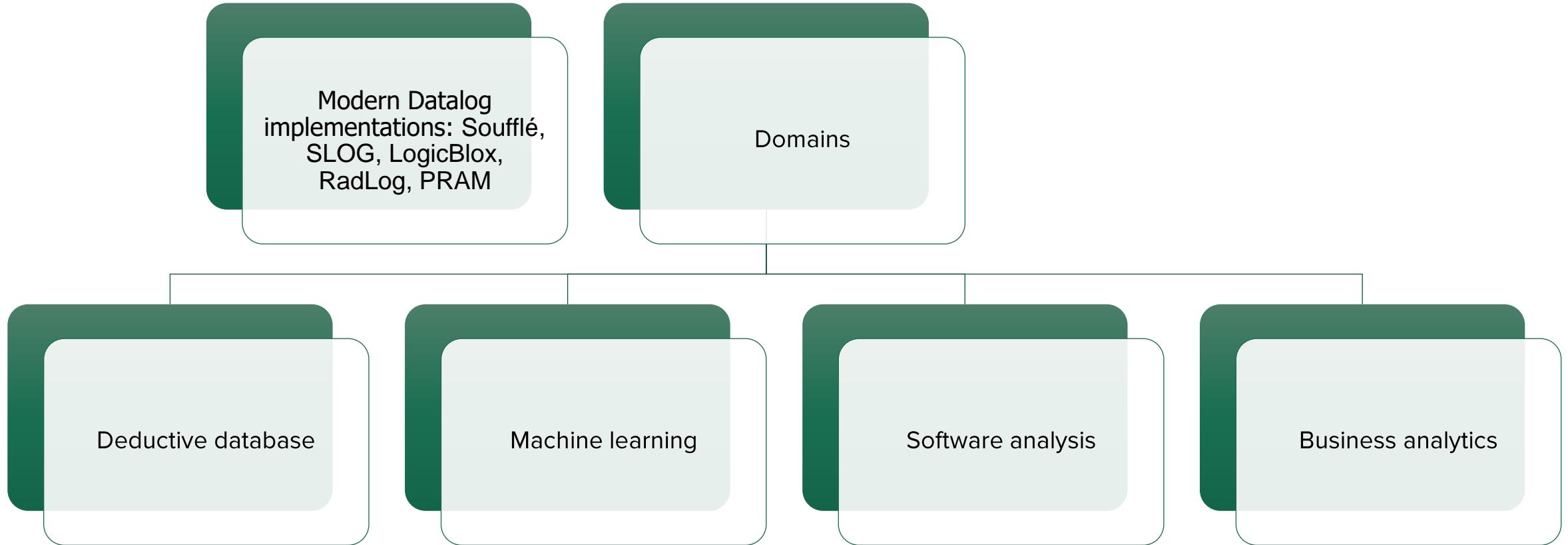
# Datalog Timeline



- Stefano Ceri, Georg Gottlob, Letizia Tanca, et al. What you always wanted to know about datalog(andnever dared to ask). *IEEE transactions on knowledge and data engineering*, 1(1):146–166, 1989.
- David Maier, K Tuncay Tekle, Michael Kifer, and David S Warren. Datalog: concepts, history, andoutlook. In *Declarative Logic Programming: Theory, Systems, and Applications*, pages 3–100. 2018.
- Shan Shan Huang, Todd Jeffrey Green, and Boon Thau Loo. Datalog and emerging applications:An interactive tutorial. In *Proceedings of the 2011 ACM SIGMOD International Conference onManagement of Data, SIGMOD ’11*, page 1213–1216, New York, NY, USA, 2011. Association forComputing Machinery.

# Datalog Applications

59



- Martin Bravenboer and Yannis Smaragdakis. Strictly declarative specification of sophisticated points-to analyses. In Proceedings of the 24th ACM SIGPLAN conference on Object oriented programming systems languages and applications, pages 243–262, 2009.
- Jiwon Seo, Stephen Guo, and Monica S Lam. Socialite: Datalog extensions for efficient social network analysis. In 2013 IEEE 29th International Conference on Data Engineering (ICDE), pages 278–289. IEEE, 2013

# Algorithm for TC computation using CUDA

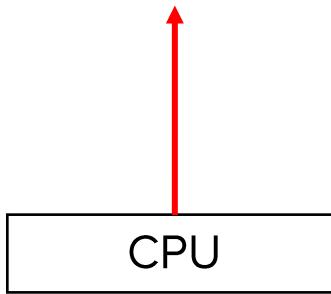
- Open-Addressing based hash table
- Two pass approach to perform hash join on the GPU
- Deduplication using sort and unique, merge and unique

```

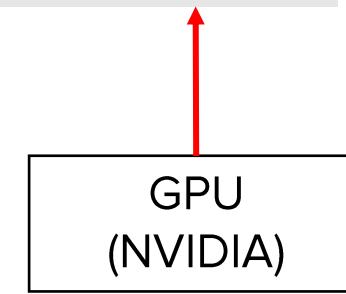
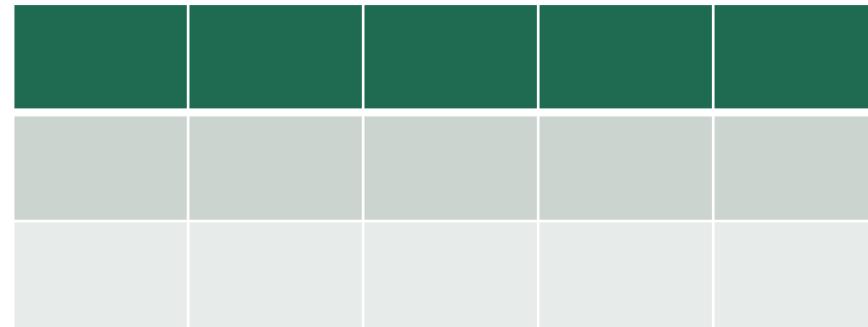
1: procedure TRANSITIVECLOSURE(Graph G)
2:    $R \leftarrow \text{HashTable}(G)$ 
3:   result  $\leftarrow \text{Sort}(G)$ 
4:    $T_\Delta \leftarrow G$ 
5:   repeat
6:     joinSizePerRow  $\leftarrow \text{JoinSize}(R, T_\Delta)$ 
7:     joinOffset  $\leftarrow \text{Scan}(\text{joinSizePerRow})$ 
8:     Initialize(joinResult, totalJoinSize)
9:     joinResult  $\leftarrow \text{Join}((R, T_\Delta), \text{joinOffset})$ 
10:    joinResult  $\leftarrow \text{Sort}(\text{joinResult})$ 
11:    joinResult  $\leftarrow \text{RemoveDuplicates}(\text{joinResult})$ 
12:    totalUniqueJoinSize  $\leftarrow \text{Size}(\text{joinResult})$ 
13:    FreeMemory( $T_\Delta$ )
14:     $T_\Delta \leftarrow \text{Copy}(\text{joinResult}, \text{totalUniqueJoinSize})$ 
15:    unionSize  $\leftarrow \text{resultSize} + \text{totalUniqueJoinSize}$ 
16:    Initialize(unionResult, unionSize)
17:    unionResult  $\leftarrow \text{MergeSortedArrays}(\text{result}, \text{joinResult})$ 
18:    unionResult  $\leftarrow \text{RemoveDuplicates}(\text{unionResult})$ 
19:    uniqueUnionSize  $\leftarrow \text{Size}(\text{unionResult})$ 
20:    oldUnionSize  $\leftarrow \text{Size}(\text{result})$ 
21:    FreeMemory(result)
22:    result  $\leftarrow \text{Copy}(\text{unionResult}, \text{uniqueUnionSize})$ 
23:    FreeMemory(joinOffset)
24:    FreeMemory(joinResult)
25:    FreeMemory(unionResult)
26:   until  $\text{oldUnionSize} \neq \text{uniqueUnionSize}$ 
27:   FreeMemory(R)
28:   FreeMemory(result)
29:   FreeMemory( $T_\Delta$ )
30:   return result
31: end procedure

```

# Off-the-shelf Data Structure



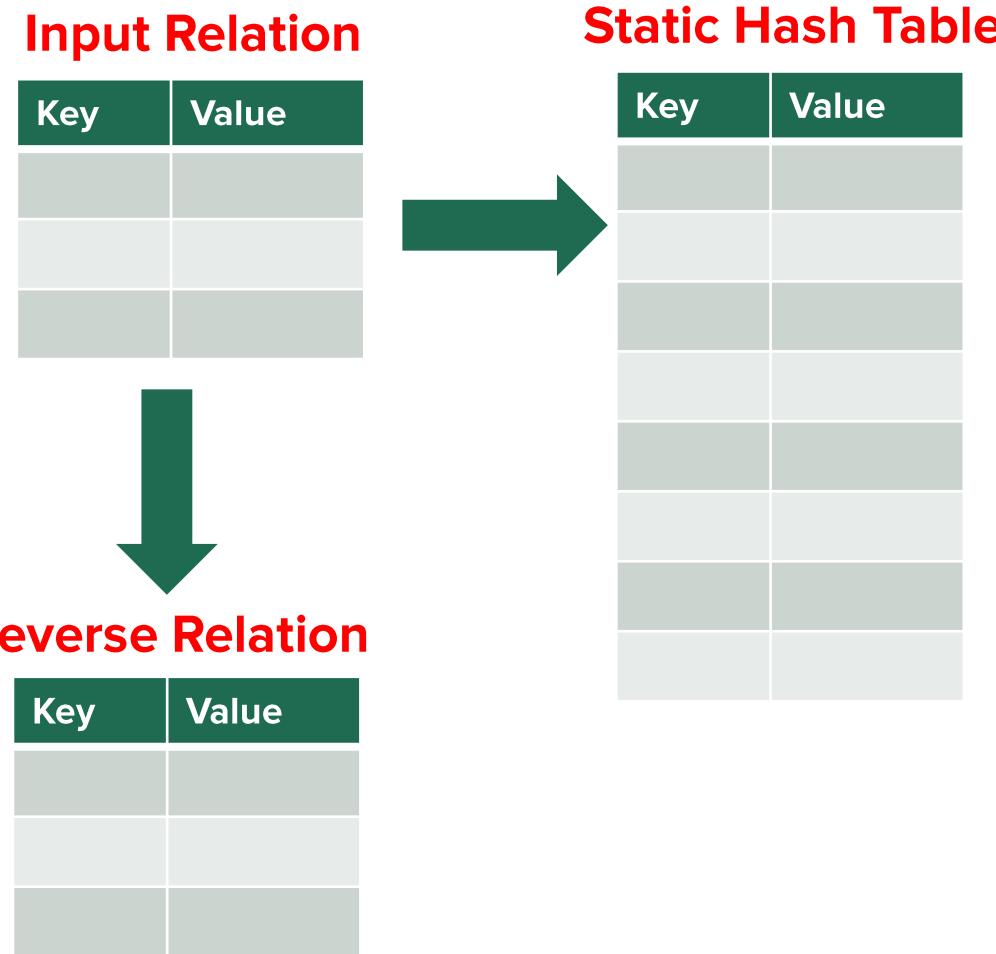
DataFrame: 2D labeled tabular data structure



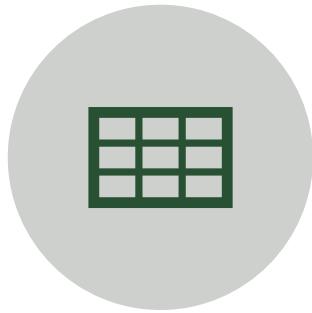
*Both supports RA primitives (e.g. join, aggregation, rename, deduplication, and projection)*

- Reback, J., McKinney, W., Van Den Bossche, J., Augspurger, T., Cloud, P., Klein, A., ... & Seabold, S. (2020). pandas-dev/pandas: Pandas 1.0. 5. Zenodo.
- Chen, D. Y. (2017). Pandas for everyone: Python data analysis. Addison-Wesley Professional.
- Green, O., Du, Z., Patel, S., Xie, Z., Liu, H., & Bader, D. A. (2021, December). Anti-Section Transitive Closure. In 2021 IEEE 28th International Conference on High Performance Computing, Data, and Analytics (HiPC) (pp. 192-201). IEEE.
- Fender, A., Rees, B., & Eaton, J. RAPIDS cuGraph. In Massive Graph Analytics (pp. 483-493). Chapman and Hall/CRC.

# Hash Join Initialization on GPU



# Why Join is Important in RA?



COMBINE DATA FROM  
MULTIPLE TABLES



FIND PATTERNS IN  
DATA



CLEAN DATA



CREATE NEW DATA  
SETS

- Daniel Zinn, Haicheng Wu, Jin Wang, Molham Aref, and Sudhakar Yalamanchili. General-purpose join algorithms for large graph triangle listing on heterogeneous systems. In Proceedings of the 9th Annual Workshop on General Purpose Processing Using Graphics Processing Unit, pages 12–21, 2016.

- A variant of Datalog for static analysis using OpenMP
- State-of-the-art implementation for multi-core CPU systems with single-node
- Translates Datalog programs to optimized C++ programs
- Supports limited number of threads for task-level parallelism
- Cannot provide data parallelism

- Herbert Jordan, Bernhard Scholz, and Pavle Subotić. Soufflé: On synthesis of program analyzers. In International Conference on Computer Aided Verification, pages 422–430. Springer, 2016.
- Thomas Gilray, Sidharth Kumar, and Kristopher Micinski. Compiling data-parallel datalog. In Proceedings of the 30th ACM SIGPLAN International Conference on Compiler Construction, CC 2021, page 23–35, New York, NY, USA, 2021. Association for Computing Machinery.

# Parallel Join (Continue)

| 65

## Design

Consider partition,  
load balancing,  
communication

## Implement

Challenging due to  
the uncertain output  
size

## Optimize

Efficient joins  
requires sorting or  
indexing

- Daniel Zinn, Haicheng Wu, Jin Wang, Molham Aref, and Sudhakar Yalamanchili. General-purpose join algorithms for large graph triangle listing on heterogeneous systems. In Proceedings of the 9th Annual Workshop on General Purpose Processing Using Graphics Processing Unit, pages 12–21, 2016.

# Hybrid Join Algorithm

- Guo et al. proposed PHYJ: SMJ with HJ join
- Reduced host-to-device and device-to-host
- Fused data communication with GPU execution
- On a single GPU achieved up to **1.72X** speedup
- Can handle skewed data
- No information on multiple GPUs or distributed systems

- Chengxin Guo, Hong Chen, Feng Zhang, and Cuiping Li. Parallel hybrid join algorithm on gpu. 2019IEEE 21st International Conference on High Performance Computing and Communications; IEEE17th International Conference on Smart City; IEEE 5th International Conference on Data Science andSystems (HPCC/SmartCity/DSS), pages 1572–1579, 2019.
- Hongzhi Wang, Ning Li, Zheng ke Wang, and Jianing Li. Gpu-based efficient join algorithms on hadoop. *The Journal of Supercomputing*, 77:292 – 321, 2020.

# Join on GPUs: Benchmark

| 67

- Rui et al. assessed NINLJ, INLJ, SMJ, and HJ on modern GPU
- Modern GPUs can lead to **20X** speedup VS **7X** speedup of old GPUs
- Not suitable for HPC systems with multiple GPU environments
- New GPU architecture is introduced (Nvidia Hopper architecture)

- Bingsheng He, Ke Yang, Rui Fang, Mian Lu, Naga Govindaraju, Qiong Luo, and Pedro Sander. Relational joins on graphics processors. In Proceedings of the 2008 ACM SIGMOD international conference on Management of data, pages 511–524, 2008.
- Ran Rui, Hao Li, and Yi-Cheng Tu. Join algorithms on gpus: A revisit after seven years. In 2015 IEEE International Conference on Big Data (Big Data), pages 2541–2550. IEEE, 2015.
- Anne C Elster and Tor A Haugdahl. Nvidia hopper gpu and grace cpu highlights. Computing in Science & Engineering, 24(2):95–100, 2022.

# Join on GPUs: LogiQL

| 68

- Wu et al. presents **Red Fox** high-performance accelerator core for **LogiQL** queries
- Outperforms multi-threaded CPU-based implementations
- Novel: multi-predicate join algorithm (worst-case optimal) on GPU
- Issue: deduplication of tuples and maintaining join result in sorted order

- Haicheng Wu, Gregory Diamos, Tim Sheard, Molham Aref, Sean Baxter, Michael Garland, and Sudhakar Yalamanchili. Red fox: An execution environment for relational query processing on gpus. InProceedings of Annual IEEE/ACM International Symposium on Code Generation and Optimization, pages 44–54, 2014.
- Haicheng Wu. Acceleration and execution of relational queries using general purpose graphics processingunit (GPGPU). PhD thesis, Georgia Institute of Technology, 2015.

# Join on GPUs: Relational Learning Framework

| 69

- Expedites rule coverage on GPUs for healthcare records data
- Outperforms **75%** of applications over multi-core CPU systems
- Duplicate tuples not efficiently managed and GPU memory overflows

- Carlos Alberto Martínez-Angeles, Haicheng Wu, Inês Dutra, Vítor Santos Costa, and Jorge BuenabadChavez. Relational learning with gpus: Accelerating rule coverage. International Journal of Parallel Programming, 44(3):663–685, 2016

# Join on GPUs: Control Flow Analysis (CFA)

| 70



Parallel functional CFA encoded in Datalog utilizes RA as the foundation on GPU



Extended Red Fox combining GPU parallelism with multi-node multi-core HPC



Proposed partitioned global address space (PGAS) programming model

- THOMAS GILRAY and SIDHARTH KUMAR. Toward parallel cfa with datalog, mpi, and cuda. InScheme and Functional Programming Workshop, 2017.

# Join on GPUs: WarpDrive

| 71

- Jünger et al. presented a single-node multi-GPU hashing for hashjoin
- Attained better memory coalescing
- Hashtable insertion rate:
  - 1.4B keys/sec (single GPU)
  - 4.3B keys/sec (4 GPUs)
- 32 bit keys only with no deduplication
- Incremental study: **WarpCore** supports 64 bit keys

- Daniel Jünger, Christian Hundt, and Bertil Schmidt. Warpdrive: Massively parallel hashing on multigpu nodes. In 2018 IEEE International Parallel and Distributed Processing Symposium (IPDPS), pages 441–450. IEEE, 2018.
- Daniel Jünger, Robin Kobus, André Müller, Christian Hundt, Kai Xu, Weigu Liu, and Bertil Schmidt. Warpcore: A library for fast hash tables on gpus. In 2020 IEEE 27th International Conference on HighPerformance Computing, Data, and Analytics (HiPC), pages 11–20, 2020.