MIPS Instruction Coding

Instruction Coding Formats

MIPS instructions are classified into four groups according to 3 coding formats, NO Pseudo-instructions or co-processor instructions required in Assignment1&2.

• R-Type - This group contains all instructions that do not require an immediate value, target offset, memory address displacement, or memory address to specify an operand. This includes arithmetic and logic with all operands in registers, shift instructions, and register direct jump instructions (jalr and jr).

All R-type instructions use opcode 000000.

• <u>I-Type</u> - This group includes instructions with an immediate operand, branch instructions, and load and store instructions. In the MIPS architecture, all memory accesses are handled by the main processor, so coprocessor load and store instructions are included in this group.

All opcodes except 000000, 00001x, and 0100xx are used for I-type instructions.

• <u>J-Type</u> - This group consists of the two direct jump instructions (j and jal). These instructions require a memory address to specify their operand.

J-type instructions use opcodes 00001x.

• <u>Coprocessor Instructions</u> - MIPS processors all have two standard coprocessors, CP0 and CP1. CP0 processes various kinds of program exceptions. CP1 is a floating point processor. The MIPS architecture makes allowance for future inclusion of two additional coprocessors, CP2 and CP3. All coprocessor instructions use opcodes 0100xx.

Note: ALL arithmetic immediate values are sign-extended. After that, they are handled as signed or unsigned 32-bit numbers, depending upon the instruction. Signed instructions can generate an overflow exception; unsigned cannot.

R-Type Instructions (Opcode 000000)

Main processor instructions that do not require a target address, immediate value, or branch displacement use an R-type

coding format. This format has fields for specifying of up to three registers and a shift amount. For instructions that do not use all of these fields, the unused fields are coded with all 0 bits. All R-type instructions use a 000000 opcode. The operation is specified by the function field.

opcode (6) rs (5) rt (5) rd (5)	sa (5) function (6)
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Alf Inc	tmustion	Function	Onad	Funct	Description	Numeric Instruction Function Funct Hex
add	rd, rs, rt		-	0x20	Add (with overflow)	Numeric instruction Function Funct flex
				0x20 $0x21$	` ´)
addu	rd, rs, rt				Add unsigned (no overflo	5W)
and	rd, rs, rt	100100	UXUU	0x24	Bitwise and	Please note that the 'break
div	rs, rt	011010	0x00	0x1A	Divide	instructions is deleted
divu	rs, rt	011011			Divide unsigned	
jalr	rd, rs	001001		0x09	Jump and link	
jr	rs rs	001001		0x08	Jump register	
mfhi	rd	010000		0x10	Move from HI	
mflo	rd	010000		0x10 $0x12$	Move from LO	
mthi	rs	010010	0x00		Move to HI	
		010001				
mtlo	rs			0x13	Move to LO	
mult	rs, rt	011000		0x18	Multiply	
multu		011001		0x19	Multiply unsigned	
nor	rd, rs, rt			0x27	Bitwise nor	
or	rd, rs, rt			0x25	Bitwise or	
sll	rd, rt, sa			0x00	Shift left logical	
sllv	rd, rt, rs	000100	0x00	0x04	Shift left logical variable	
slt	rd, rs, rt	101010	0x00	0x2A	Set on less than (signed)	
sltu	rd, rs, rt	101011	0x00	0x2B	Set on less than unsigned	
sra	rd, rt, sa	000011	0x00	0x03	Shift right arithmetic	
srav	rd, rt, rs	000111	0x00	0x07	Shift right arithmetic vari	iable
srl	rd, rt, sa	000010	0x00	0x02	Shift right logical	
srlv	rd, rt, rs	000110	0x00	0x06	Shift right logical variable	le
sub	rd, rs, rt	100010	0x00	0x22	Subtract	
subu	rd, rs, rt	100011	0x00	0x23	Subtract unsigned	
syscall		001100	0x00	0x0C	System call	
xor	rd, rs, rt	100110	0x00	0x26	Bitwise exclusive or	

I-Type Instructions (All opcodes except 000000, 00001x, and 0100xx)

I-type instructions have a 16-bit immediate field that codes an immediate operand, a branch target offset, or a displacement for a memory operand. For a branch target offset, the immediate field contains the signed difference between the address of the following instruction and the target label, with the two low order bits dropped. The dropped bits are always 0 since instructions are word-aligned.

immediate (16)

For the bgez, bgtz, blez, and bltz instructions, the rt field is used as an extension of the opcode field.

rt (5)

	code (0)	15 (5)		(5)		miniediate (10)
Alf Ir	struction	Opcode	Notes	Oncd	Des	scription	
	rt, rs, immediate	-		-		d immediate (with overflow)	
	rt, rs, immediate			0x09	Ado	d immediate unsigned (no overflow)	
andi	rt, rs, immediate	001100		0x0C	Bitv	wise and immediate	
beq	rs, rt, label	000100		0x04	Bra	nch on equal	
bgez	rs, label	000001	rt=00001	0x01	Bra	nch on greater than or equal to zero	
bgtz	rs, label	000111	rt=00000	0x07	Bra	nch on greater than zero	
blez	rs, label	000110	rt=00000	0x06	Bra	nch on less than or equal to zero	
bltz	rs, label	000001	rt=00000	0x01	Bra	nch on less than zero	
bne	rs, rt, label	000101		0x05	Bra	nch on not equal	
lb	rt, immediate(rs)	100000		0x20	Loa	ad byte	
lbu	rt, immediate(rs)	100100		0x24	Loa	nd byte unsigned	
lh	rt, immediate(rs)	100001		0x21	Loa	nd halfword	
lhu	rt, immediate(rs)	100101		0x25	Loa	nd halfword unsigned	
lui	rt, immediate	001111		0x0F	Loa	ad upper immediate	
lw	rt, immediate(rs)	100011		0x23	Loa	nd word	
ori	rt, rs, immediate	001101		0x0D) Bit	twise or immediate	
sb	rt, immediate(rs)	101000		0x28	3 Sto	re byte	
slti	rt, rs, immediate	001010		0x0A	A Set	t on less than immediate (signed)	
sltiu	rt, rs, immediate	001011		0x0B	3 Set	on less than immediate unsigned	
sh	rt, immediate(rs)	101001		0x29	Sto	re halfword	
sw	rt, immediate(rs)	101011		0x2B	3 Sto	ore word	
xori	rt, rs, immediate	001110		0x0E	E Bit	wise exclusive or immediate	
lwl	rt, immediate(rs)	100010		0x22	2 Lo	ad word left (unaligned)	Please note that the last 4
lwr	rt, immediate(rs)	100110		0x26	5 Lo	ad word right (unaligned)	instructions are updated
swl	rt, immediate(rs)	101010		0x2A	A Sto	ore word left (unaligned)	(lwl, lwr, swl, swr)
		101110		0.01	- a.	1 1 1 / 1 1	

J-Type Instructions (Opcode 00001x)

rt, immediate(rs) 101110

swr

opcode (6)

rs (5)

The only J-type instructions are the jump instructions j and jal. These intructions require a 26-bit coded address field to specify the target of the jump. The coded address is formed from the bits at positions 27 to 2 in the binary representation of the address. The bits at positions 1 and 0 are always 0 since instructions are word-aligned.

0x2E Store word right (unaligned)

When a J-type instruction is executed, a full 32-bit jump target address is formed by concatenating the high order four bits of the PC (the address of the instruction following the jump), the 26 bits of the target field, and two 0 bits.

1 (6)	torget (26)
opcode (6)	target (26)

Instruction Opcode Target

Opcd Description

j label 000010 coded address of label 0x02 Jump

jal label 000011 coded address of label 0x03 Jump and link

Syscalls you need to support in Project2(not required in project1)

The syscalls you need to support are: **1, 4, 5, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17** in the following chart.

Service	System call code	Arguments	Result	
print_int	1	\$a0 = integer		
print_float	2	\$f12 = float		
print_double	3	\$f12 = double		
print_string	4	\$a0 = string		
read_int	5		integer (in \$v0)	
read_float	6		float (in \$f0)	
read_double	7		double (in \$f0)	
read_string	8	\$a0 = buffer, \$a1 = length		
sbrk	9	\$a0 = amount	address (in \$v0)	
exit	10			
print_char	11	\$a0 = char		
read_char	12		char (in \$v0)	
open	13	\$a0 = filename (string), \$a1 = flags, \$a2 = mode	file descriptor (in \$a0)	
read	14	\$a0 = file descriptor, \$a1 = buffer, \$a2 = length	num chars read (in \$a0)	
write	15	\$a0 = file descriptor, \$a1 = buffer, \$a2 = length	num chars written (in \$a0)	
close	16	\$a0 = file descriptor		
exit2	17	\$a0 = result		

Page URL: http://www.cs.sunysb.edu/~lw/spim/MIPSinstHex.pdf from http://www.d.umn.edu/~gshute/spimsal/talref.html

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How syscall works?

Recall that you will read in a MIPS file as input in your project 1, and you will need to run the code (Read Project 1 instruction first if you do not understand). In your assembling part, you simply need to put the binary code: 0000000000000000000000001100 in, all syscalls have the same machine code.

How do we distinguish them then? Syscalls are distinguished by checking the value stored in \$v0 register. In your simulation part, when ever you see a syscall made (seeing this binary code: 0000000000000000000000000001100), you simply go and check what is stored in \$v0. In other words, you can write a switch, with the value in \$v0 being the cases. For each case, you simply implement its functionality. Using print_int as an example, in your case: 1, you can have printf("%d", *a0); . The argument column specifies the arguments this syscall takes (the integer to be printed is stored in \$a0, for example).

The input MIPS code format

You will need to consider the following situations while reading the input MIPS file:

- 1. There will only be .data and .text sections.
- 2. There could be spaces or tabs before and after each line.
- 3. There could be spaces before and after each element within a line. e.g. add \$t0, \$t1, \$t2.
- 4. There could be empty lines.
- 5. There could be comments after the line of code. There could also be a line with only comments. Comments are always following a "#".
- 6. Labels can be followed by a line of code, or can have it's own line. Labels are labeling the same line of code in both situations.

```
case1
label: add $t0, $t1, $t2

case2
label:
add $t0, $t1, $t2
```

The data types you need to support (for project 2)

The data types you need to support are:

- 1. ascii
- 2. asciiz
- 3. word
- 4. byte
- 5. half