Report

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Selected Paper: Te-Yuan Huang, Ramesh Johari, Nick McKeown. **Downton Abbey Without the Hiccups: Buffer-Based Rate Adaptation for HTTP Video Streaming**.

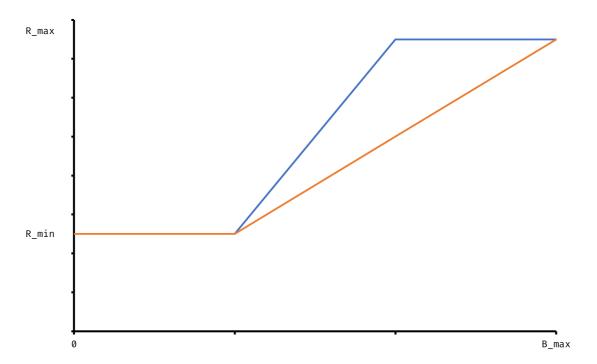
Algorithm Analysis

Algorithm: Video Rate Adaptation Algorithm

```
1
    Hyperparameters:
2
        R_M: All available bitrates, bits per second
 3
        R_max: Maximum available bitrate, bits per second
4
        R_min: Minimum available bitrate, bits per second
 5
        R_i: i-th available bitrate inside R_M, bits per second
        mapping_func: a mapping function to transfer buffer occupancy to
 6
    designated bitrate mathematically.
7
8
    Input:
9
        Rate_prev: The previously used video rate, bits per second
10
        Buf_now: The current buffer occupancy, second
11
12
    Output:
13
        Rate_next: The fetched video rate, second
14
15
    # Select adjacent possible values from R_M, larger and smaller respectively
   if Rate_prev == R_max:
16
        Rate_plus = R_max # Could not increase any more
17
18
    else:
19
        Rate_plus = min{R_i : R_i > Rate_prev} # next larger step
   if Rate_prev == R_min:
20
        Rate_minus = R_min # Could not decrease any more
21
22
    else:
23
        Rate_minus = max{R_i : R_i < Rateprev} # next smaller step</pre>
24
    # Check the mapping function to see whether should move up/down to next
25
    step. If so, move to the step nearest to the bound, skipping intervals (if
    exists).
    if mapping_func(Buf_now) >= Rate_plus:
26
        Rate_next = max{R_i : R_i < mapping_func(Buf_now)}</pre>
27
28
    elif mapping_func(Buf_now) <= Rate_minus:</pre>
29
        Rate_next = min{R_i : R_i > mapping_func(Buf_now)}
30
    else
31
        Rate_next = Rate_prev
32
33
    return Ratenext
```

Mapping Function

The orange line represents the naive implementation from the paper, while the blue line is my enhanced version. By selecting the R_max at lower buffer occupancy, a.k.a. B_t, we are able to achieve higher average bitrate. According to the paper, the plain part for small B_t, reservoir, is with the length of $V \times \frac{R_{\max}}{R_{\min}}$, where V is the duration of a chunk.



I set the upper plain part with the length of **twice** as **reservoir** if possible. If this is not a valid value, it will try to find the most nearest valid one, moving to <code>B_max</code> one by one. This enhanced mapping function gives better performance, especially on good network condition (faster convergence).

Implementation

```
def mapping_func(buffer_now): 4 usages
    """
    The mapping function between buffer occupancy and bitrate.

Args:
    buffer_now: The current buffer occupancy, in seconds

Returns:
    The corresponding bitrate
    """

global RESERVOIR, B_MAX, R_MIN, R_MAX, RMAX_POINT

if buffer_now 
RESERVOIR:
    return R_MIN

elif buffer_now 
RMAX_POINT:
    return R_MAX

else:
    return R_MIN + (R_MAX - R_MIN) * (buffer_now - RESERVOIR) / (RMAX_POINT - RESERVOIR)
```

```
def video_rate_adaptation_algo(rate_prev, buffer_now): 1 usage
   This function is called every time a new video chunk is requested.
       rate_prev: The previously used video rate, in bits per second
       buffer_now: The current buffer occupancy, in seconds
   Returns:
       The next video rate
   global R_MIN, R_MAX, AVAILABLE_BITRATES
   if rate_prev = 0:
       return R_MIN # Fast-track the case when rate_prev is 0 for the first chunk
   if rate_prev = R_MAX:
       rate_plus = R_MAX
       rate_plus = min([rate for rate in AVAILABLE_BITRATES if rate ≥ rate_prev])
   if\ rate\_prev\ =\ R\_MIN: # Consider the case when rate\_prev is 0 for the first chunk
       rate_minus = R_MIN
       rate_minus = max([rate for rate in AVAILABLE_BITRATES if rate ≤ rate_prev])
   if mapping_func(buffer_now) ≥ rate_plus:
       rate_next = max([rate for rate in AVAILABLE_BITRATES if rate ≤ mapping_func(buffer_now)])
   elif mapping_func(buffer_now) ≤ rate_minus:
       rate_next = min([rate for rate in AVAILABLE_BITRATES if rate ≥ mapping_func(buffer_now)])
       rate_next = rate_prev
   return rate_next
```

Evaluation

Naive / Mine	Avg. Bitrate	Buffer Time	Switches	Score
badtest	2166666 / 633333	73 / 1	23 / 1	7518 / 553192
testALThard	2166666 / 633333	72 / 1	22 / 1	8535 / 553504
testALTsoft	3816666 / 1583333	27 / 0.202	10 / 2	407852 / 1326319
testHD	4566666 / 1583333	0.202 / 0.202	2/2	3825384 / 1326319
testHDmanPQtrace	50000 / 50000	242 / 242	0/0	0.197 / 0.197
testPQ	50000 / 50000	246 / 246	0/0	0.162 / 0.162

This is the comparison table between the example code and my implementation. We can easily find out the last two row share the same statistics, which means that they perform similarly under bad network condition. As for the first three ones, they indicate that my implementation outperform the given one on fluctuated network. For the HD test, it shows that my implementation does not achieve the highest bitrate as quick as the given one.