

(10.14) Poisson equation in 2D on unit square

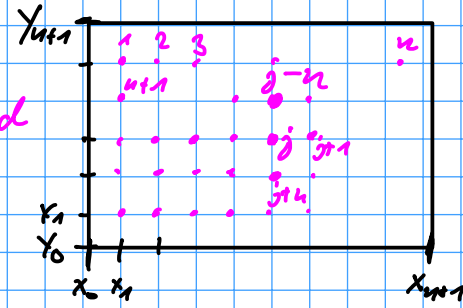
$$\Delta u = f, \quad u|_{\partial\Omega} = 0$$

discretize Δ :

use 1D stencil in every direction

$$\frac{1}{h^2} \begin{bmatrix} & & -1 \\ -1 & & \\ & 2+2 & -1 \\ & & -1 \end{bmatrix}$$

Ω_h
discrete grid



start with
 Ω unit
square

use geometry
of Ω

$$= \frac{1}{h^2} \begin{bmatrix} & & -1 \\ -1 & 4 & -1 \\ & & -1 \end{bmatrix}$$

5-point-stencil

apply stencil at every unknown grid point
 $\Rightarrow n^2$ equations

unknowns = $n \cdot n = n^2$

vector of unknowns

$$u_h = \begin{pmatrix} u_1 \\ u_2 \\ \vdots \\ u_{n^2} \end{pmatrix}$$

j th equation:

$$\frac{1}{h^2} (0 \dots 0 \underset{\substack{\uparrow \\ (j-n)\text{th} \\ \text{col.}}}{-1} 0 \dots 0 \underset{\substack{\uparrow \\ j\text{th} \\ \text{col.}}}{-1} 4 \underset{\substack{\uparrow \\ (j+n)\text{th} \\ \text{col.}}}{-1} 0 \dots 0 -1 0 0) \begin{pmatrix} u_1 \\ u_2 \\ \vdots \\ u_{n^2} \end{pmatrix} = f_j$$

if this is a grid point that has 4 neighbors that are unknowns?
(inner grid point)

if neighbors are at the boundary then the value of this neighbor is known.
 These values go to the right-hand side of the equation.

e.g. no left neighbor

$$\frac{1}{h^2} (-1 \dots -1 \quad \textcircled{4} -1 \dots -1) u_h = f_j + \underbrace{\frac{1}{h^2} u}_{=0}(\text{boundary})$$

because of $u|_{\partial\Omega} = 0$

the whole system ($n^2 \times n^2$ system):

$$\frac{1}{h^2} \begin{pmatrix} \begin{matrix} 4 & -1 & & \\ -1 & 4 & -1 & \\ & -1 & 4 & -1 \\ & & -1 & 4 \end{matrix} & \begin{matrix} -1 & & & \\ & -1 & & \\ & & -1 & \\ & & & -1 \end{matrix} \\ \begin{matrix} -1 & & & \\ & -1 & & \\ & & -1 & \\ & & & -1 \end{matrix} & \begin{matrix} 4 & -1 & & \\ -1 & 4 & -1 & \\ & -1 & 4 & -1 \\ & & -1 & 4 \end{matrix} & \begin{matrix} -1 & & & \\ & -1 & & \\ & & -1 & \\ & & & -1 \end{matrix} \\ \vdots & \vdots & \vdots & \vdots \end{pmatrix} \begin{pmatrix} u_1 \\ u_2 \\ \vdots \\ u_{n^2} \end{pmatrix} = \begin{pmatrix} f_1 \\ \vdots \\ f_{n^2} \end{pmatrix}$$

$+0_L + 0_T$
 $+0_E$
 \vdots
 $+0_S$
 $+0_E + 0_R$

$$T_h u_h = f_h \quad (n^2 \times n^2) \quad (n^2 \times 1)$$

10.15 Compact writing: Kronecker products

1D case: $\begin{pmatrix} 2 & -1 & & \\ -1 & 2 & -1 & \\ & -1 & \ddots & -1 \\ & & -1 & 2 \end{pmatrix} = A_n^{(1)}$

$n=3 \leadsto u_1, \dots, u_9 \rightarrow$ 1D discrete Laplace + $2I$

$$\frac{1}{h^2} \begin{pmatrix} \boxed{\begin{matrix} 4 & -1 \\ -1 & 4 & -1 \\ -1 & 4 \end{matrix}} & \begin{matrix} -1 \\ -1 \\ -1 \end{matrix} & \\ \begin{matrix} -1 \\ -1 \\ -1 \end{matrix} & \begin{matrix} 4 & -1 \\ -1 & 4 & -1 \\ -1 & -1 & 4 \end{matrix} & \begin{matrix} -1 \\ -1 \\ -1 \end{matrix} \\ & \begin{matrix} -1 \\ -1 \\ -1 \end{matrix} & \begin{matrix} 4 & -1 \\ -1 & 4 & -1 \\ -1 & -1 & 4 \end{matrix} \end{pmatrix}$$

$A_n^{(2)}$

$$= \frac{1}{h^2} \begin{pmatrix} A_n^{(1)} + 2I & -I & \\ -I & A_n^{(1)} + 2I & -I \\ & -I & A_n^{(1)} + 2I \end{pmatrix}$$

$$= \frac{1}{h^2} \left(\begin{pmatrix} A_n^{(1)} & & & \\ & A_n^{(1)} & & \\ & & A_n^{(1)} & \\ & & & A_n^{(1)} \end{pmatrix} + \begin{pmatrix} 2I & -I & & \\ -I & 2I & -I & \\ & -I & 2I & \\ & & -I & 2I \end{pmatrix} \right)$$

looks like $A_n^{(1)}$ with all numbers replace by mult. of I

Kronecker Tensor product of two matrices A, B :

$$A \otimes B = \begin{pmatrix} a_{11} B & \dots & a_{1n} B \\ \vdots & & \vdots \\ a_{m1} B & \dots & a_{mn} B \end{pmatrix}$$

in Matlab

`kron(A, B)`

works efficiently for sparse matrices

$$A_n^{(2)} = \underbrace{I \otimes A_n^{(1)}}_{\begin{pmatrix} 1 \cdot A_n^{(1)} & & \\ & 1 \cdot A_n^{(1)} & \\ & & \ddots \\ & & & 1 \cdot A_n^{(1)} \end{pmatrix}} + \underbrace{A_n^{(1)} \otimes I}_{\begin{pmatrix} 2 \cdot I & 1 \cdot I & & \\ -1 \cdot I & 2 \cdot I & & \\ & \ddots & \ddots & \\ -1 \cdot I & 2 \cdot I & & \end{pmatrix}}$$

in Matlab:

$$A2 = \text{kron}(\text{eye}(n), A1)$$

$$+ \text{kron}(A1, \text{eye}(n))$$

`eye(n)` sparse identity matrix

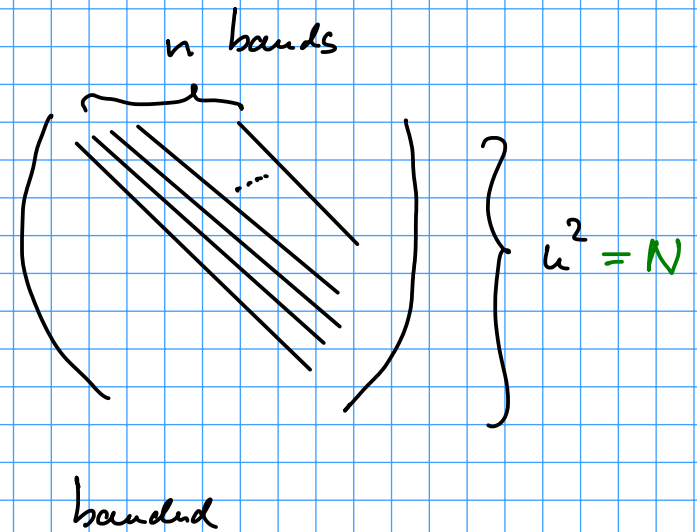
Compare

$$\Delta_{\text{in 2D}} = \frac{\partial^2}{\partial x^2} + \frac{\partial^2}{\partial y^2}$$

(10.16) Complexity of Cholesky factorization

Observation: fill-in between inner and outer diagonals:

$R = \text{chol}(A)$ has the structure



full Cholesky would need $O(N^3) = O(u^6)$

banded Cholesky needs $O(N^2)$

$$N^2 = \underbrace{u^2}_{\text{\#rows}} \cdot \underbrace{u^2}_{\text{bandwidth}^2}$$

concrete experiments:

$$u = 1000$$

$$\text{CPU-time} \approx 15 \text{ s}$$

$$u = 2000$$

$$\text{CPU-time} \approx 240 \text{ s}$$

ultimate goal would be $O(N)$ computing time

$O(N)$ is possible with multigrid methods

but $O(N \log N)$ is possible for Poisson using FFT \longrightarrow next week

